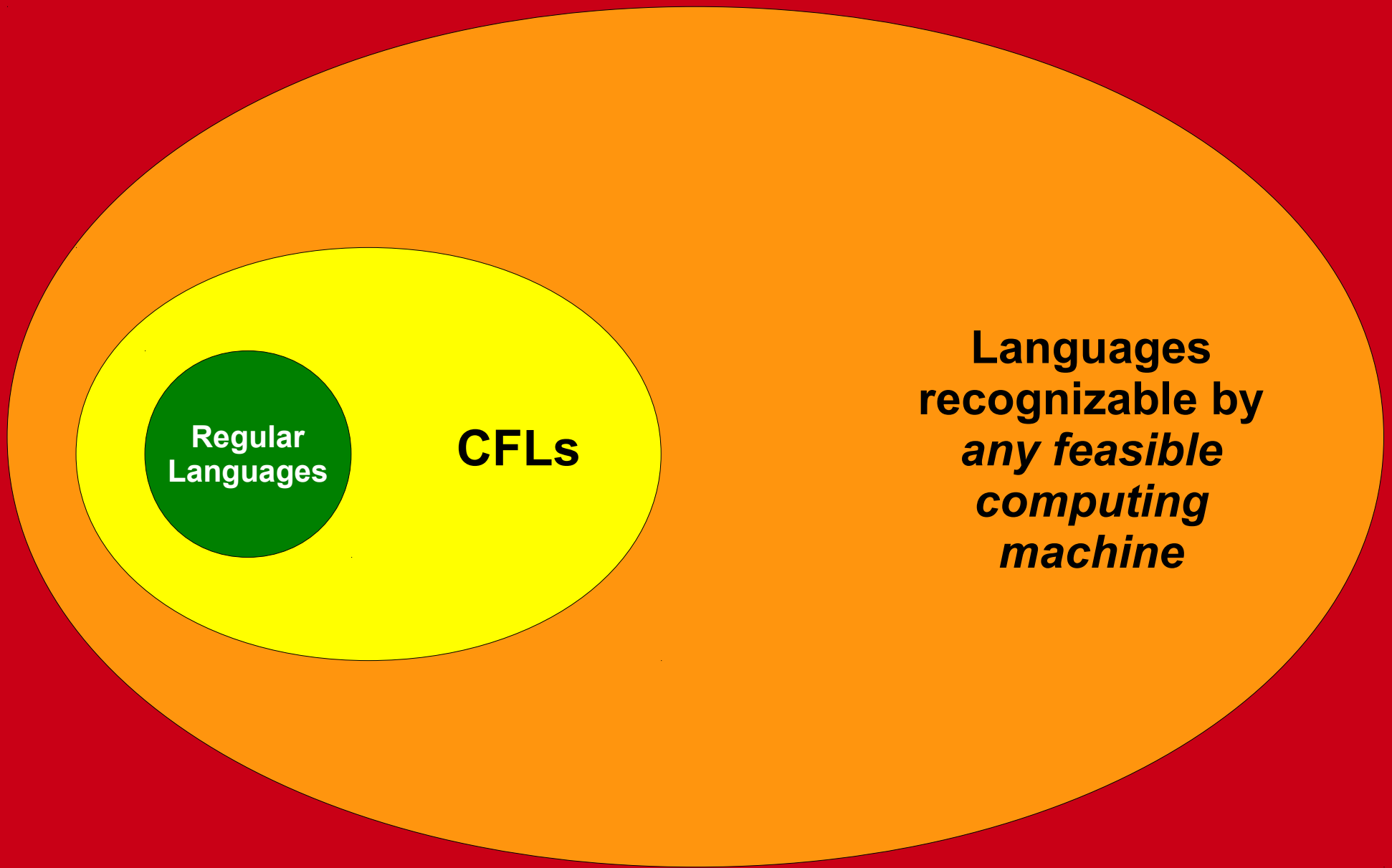


# Turing Machines

## Part One

What problems can we solve with a computer?



Regular Languages

CFLs

Languages recognizable by *any feasible computing machine*

All Languages

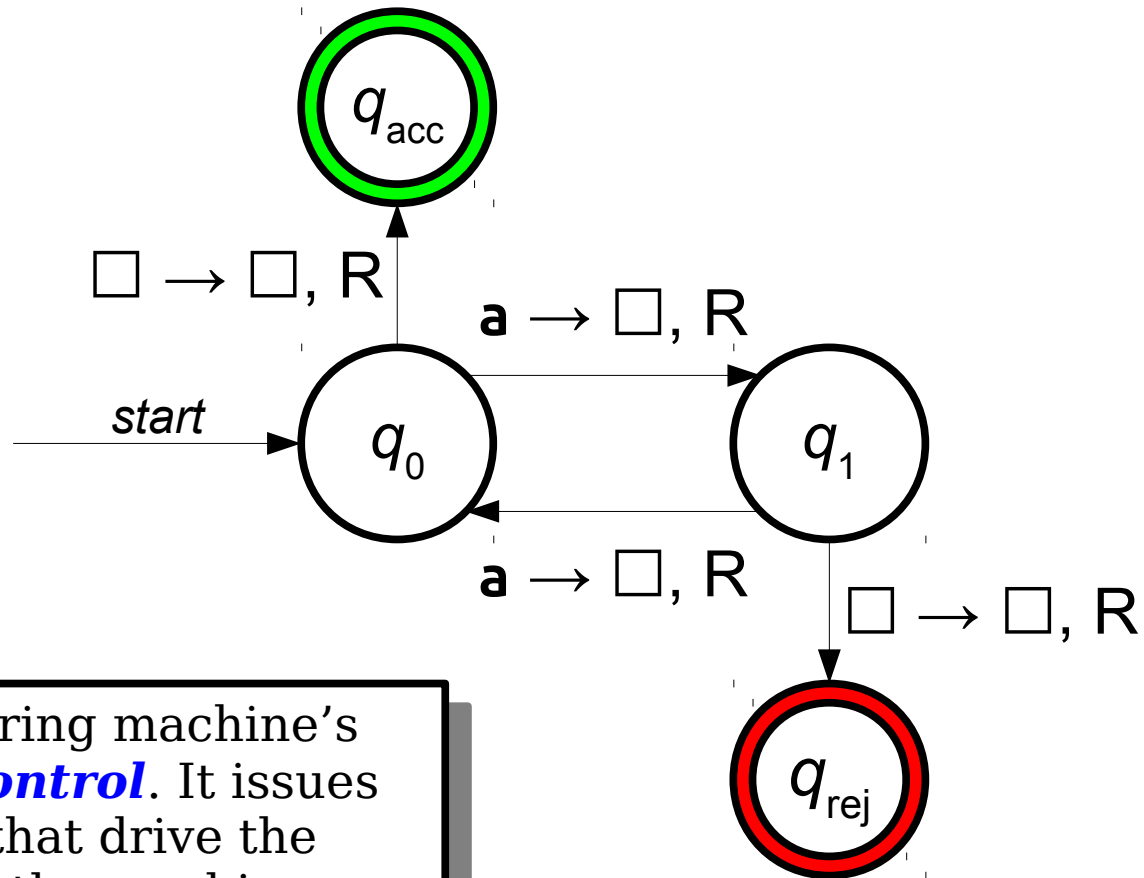
That same drawing, to scale.

# The Problem

- Finite automata accept precisely the regular languages.
- We may need unbounded memory to recognize context-free languages.
  - e.g.  $\{ a^n b^n \mid n \in \mathbb{N} \}$  requires unbounded counting.
- How do we build an automaton with finitely many states but unbounded memory?

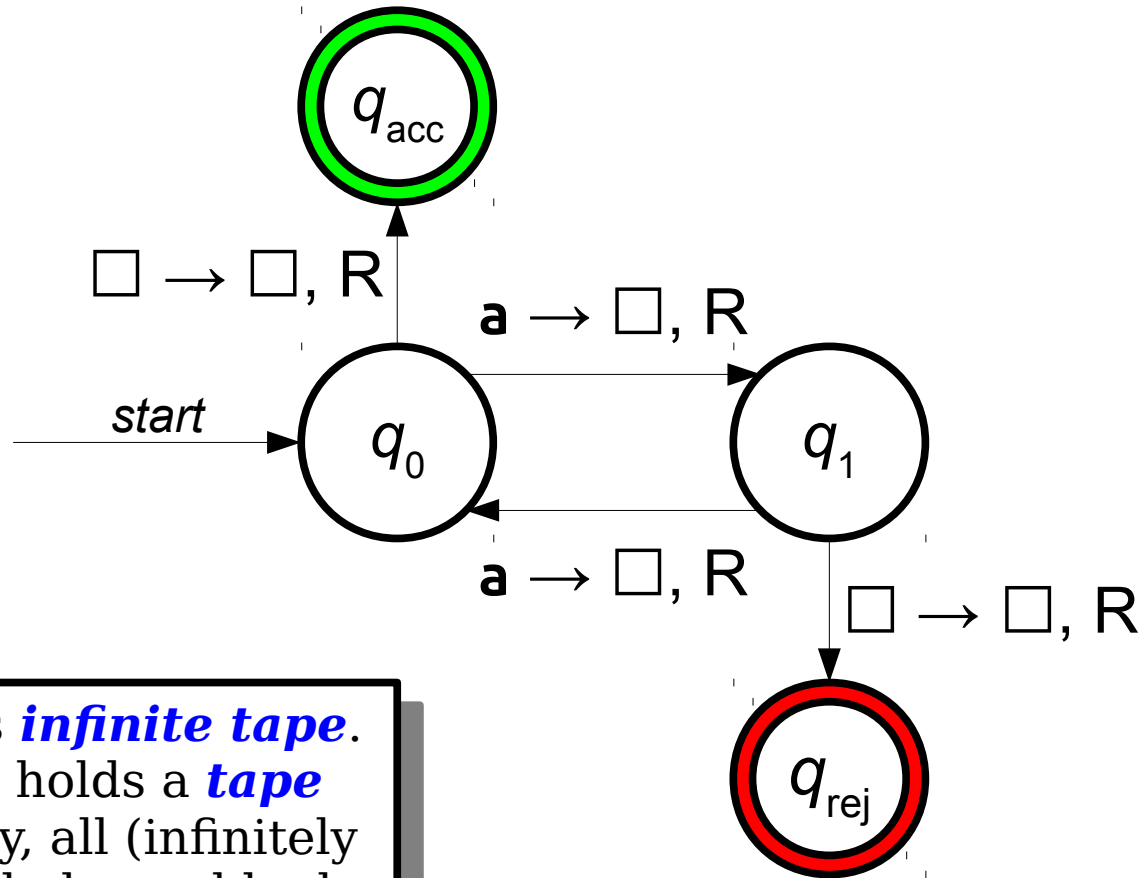
# A Brief History Lesson

# A Simple Turing Machine

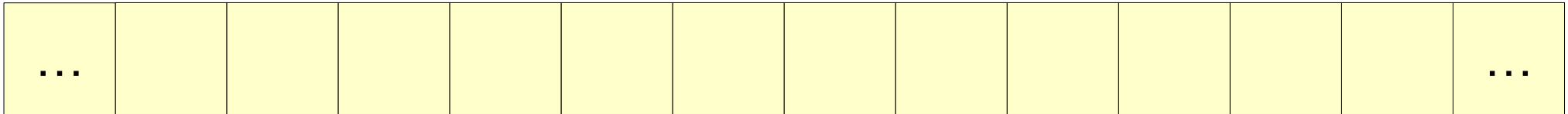


This is the Turing machine's **finite state control**. It issues commands that drive the operation of the machine.

# A Simple Turing Machine

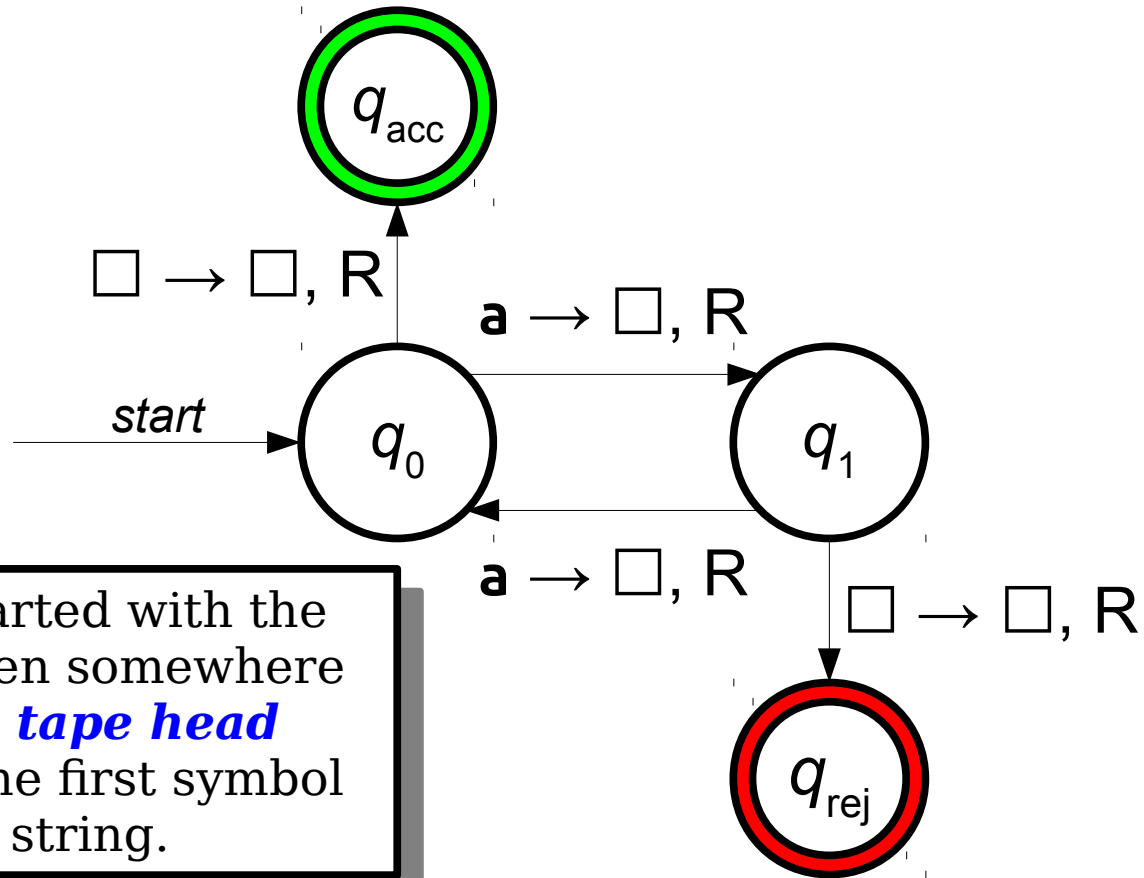


This is the TM's *infinite tape*. Each tape cell holds a *tape symbol*. Initially, all (infinitely many) tape symbols are blank.

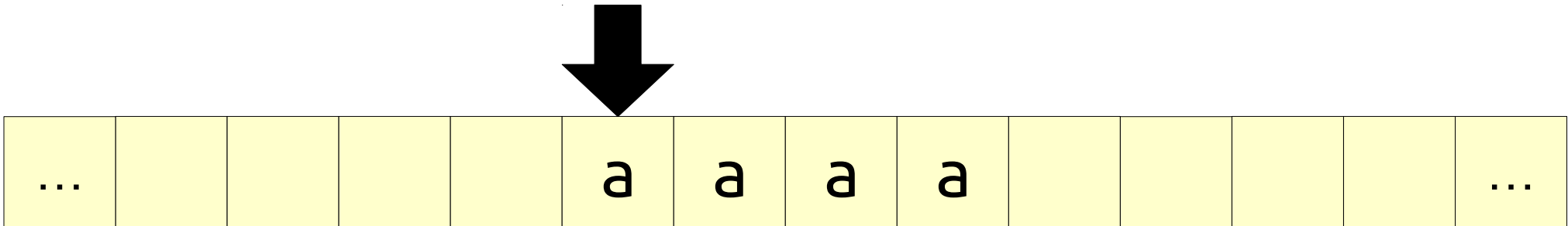




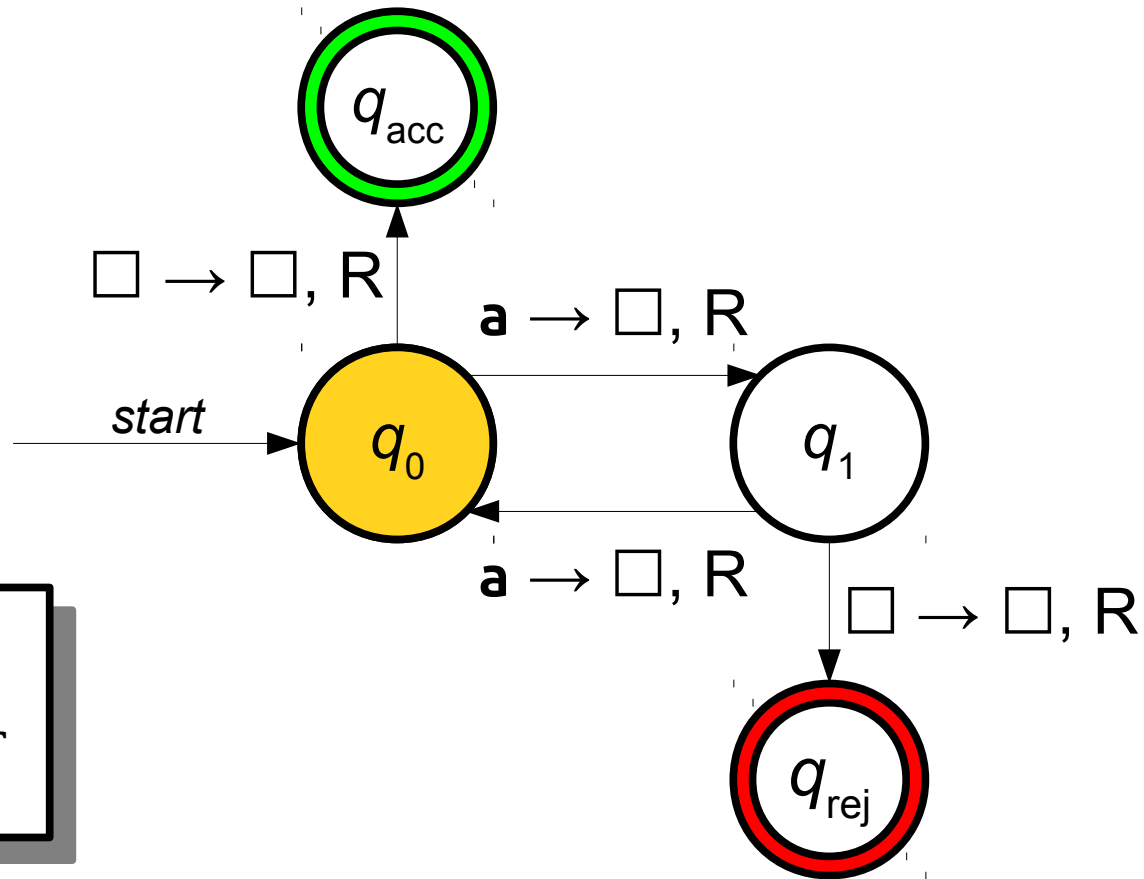
# A Simple Turing Machine



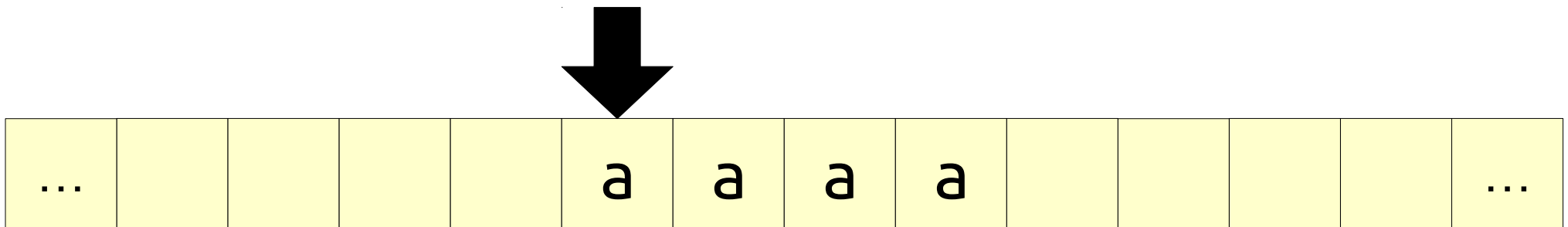
The machine is started with the **input string** written somewhere on the tape. The **tape head** initially points to the first symbol of the input string.



# A Simple Turing Machine

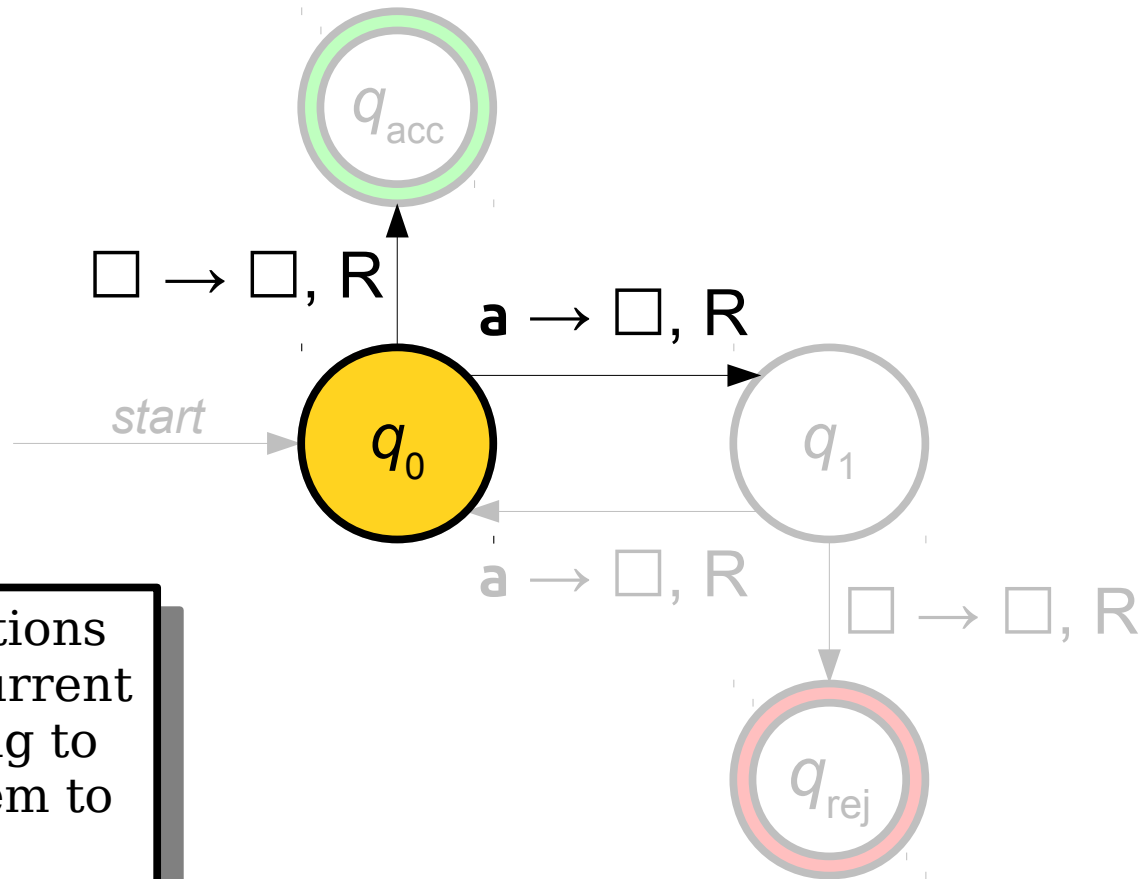


Like DFAs and NFAs, TMs begin execution in their **start state**.

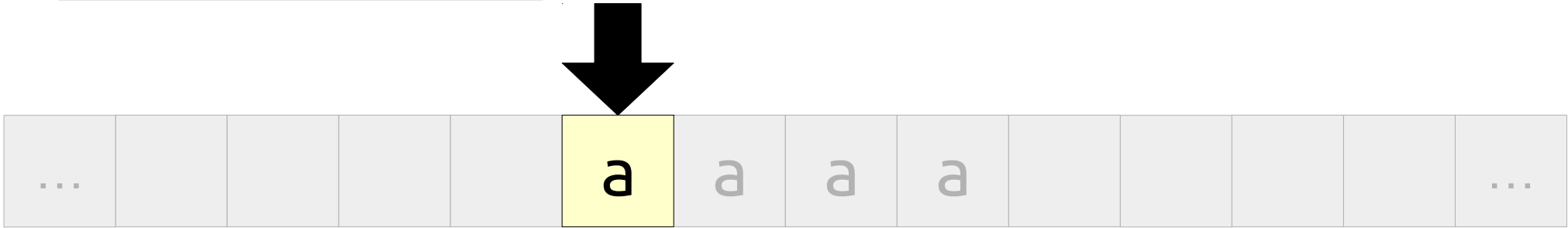




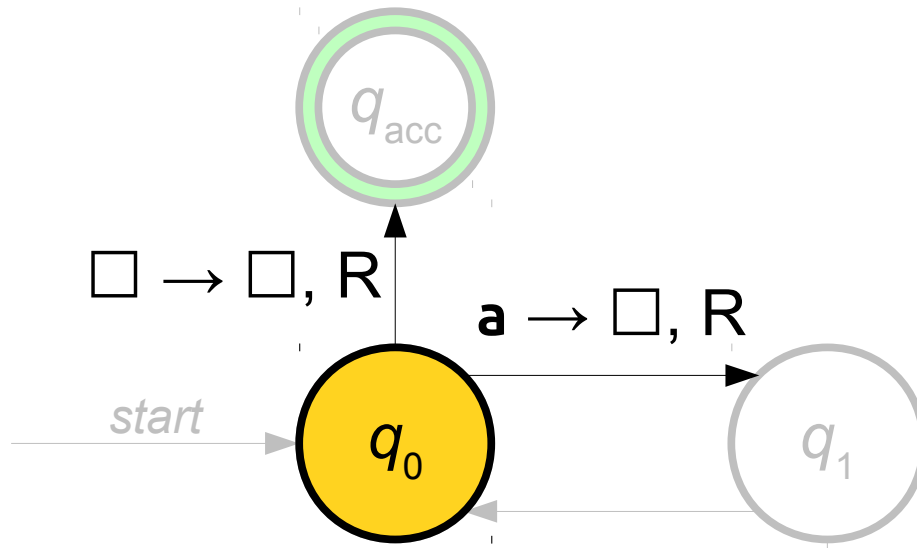
# A Simple Turing Machine



These two transitions originate at the current state. We're going to choose one of them to follow.



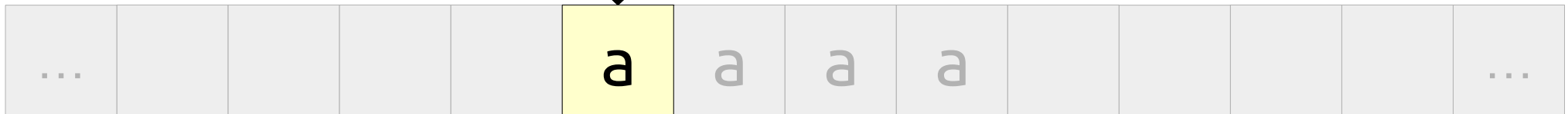
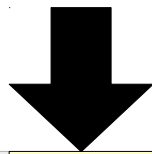
# A Simple Turing Machine



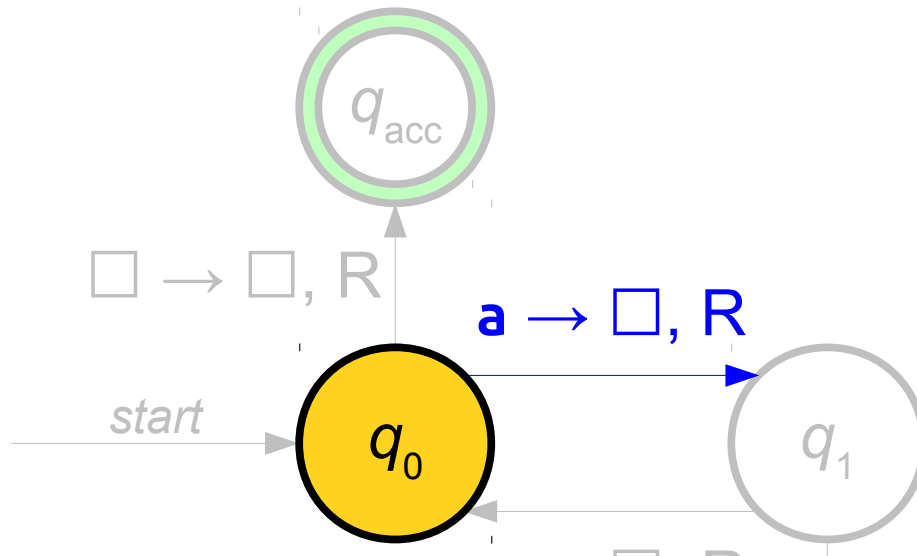
Each transition has the form

***read*** → ***write***, ***dir***

and means “if symbol ***read*** is under the tape head, replace it with ***write*** and move the tape head in direction ***dir*** (L or R). The □ symbol denotes a blank cell.



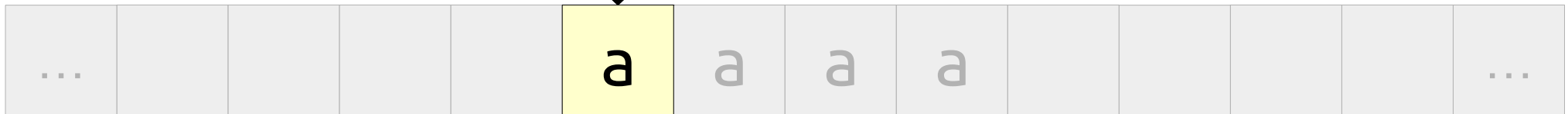
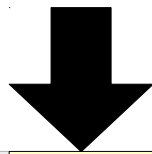
# A Simple Turing Machine



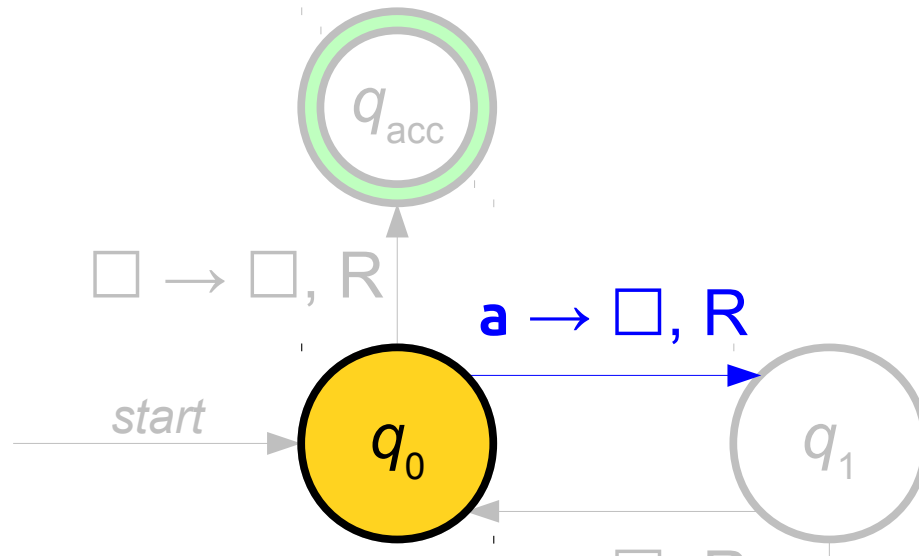
Each transition has the form

***read*** → ***write***, ***dir***

and means “if symbol ***read*** is under the tape head, replace it with ***write*** and move the tape head in direction ***dir*** (L or R). The □ symbol denotes a blank cell.



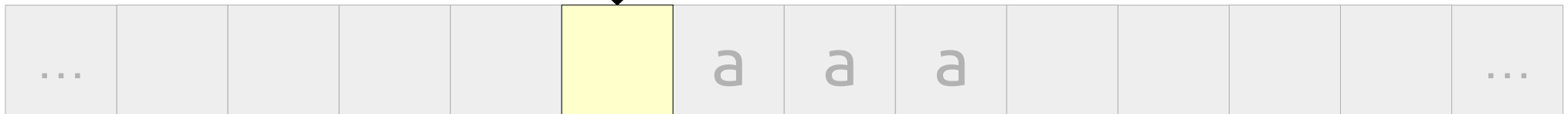
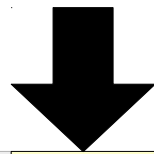
# A Simple Turing Machine



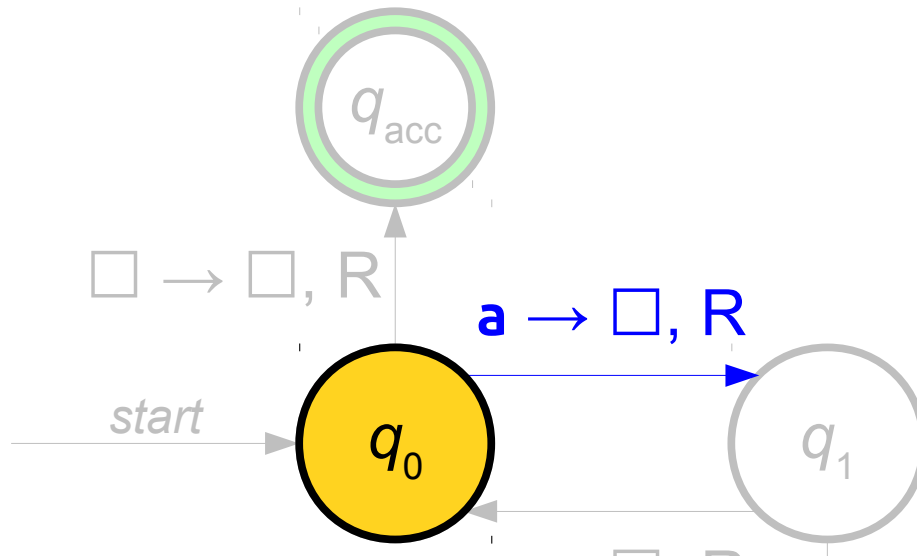
Each transition has the form

***read*** → ***write***, ***dir***

and means “if symbol ***read*** is under the tape head, replace it with ***write*** and move the tape head in direction ***dir*** (L or R). The □ symbol denotes a blank cell.



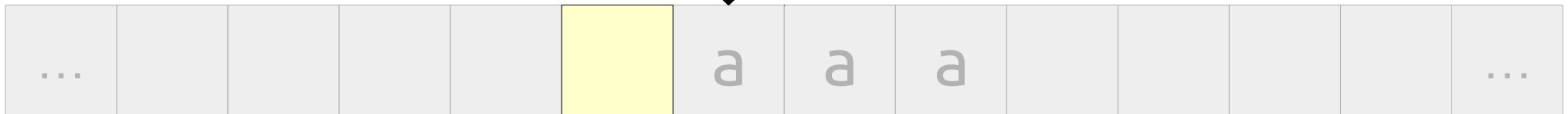
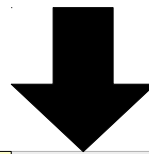
# A Simple Turing Machine



Each transition has the form

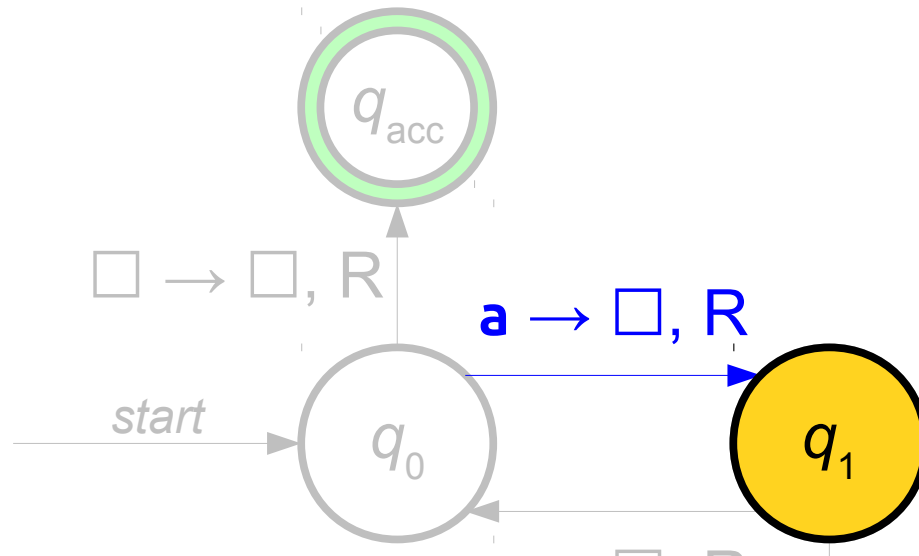
***read*** → ***write***, ***dir***

and means “if symbol ***read*** is under the tape head, replace it with ***write*** and move the tape head in direction ***dir*** (L or R). The □ symbol denotes a blank cell.





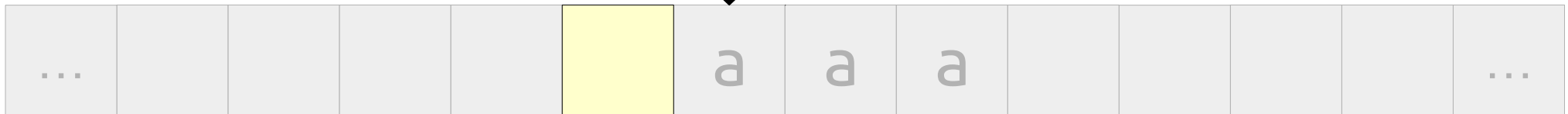
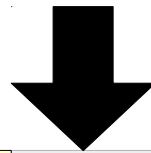
# A Simple Turing Machine



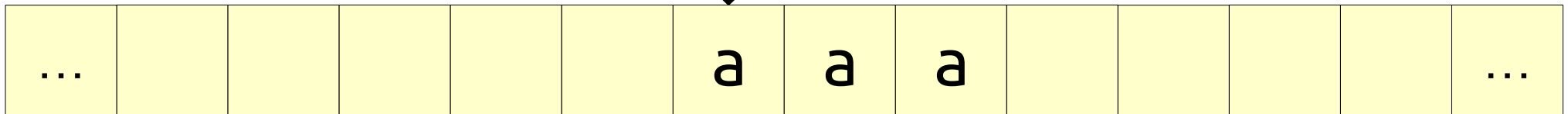
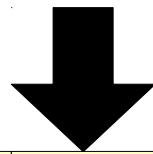
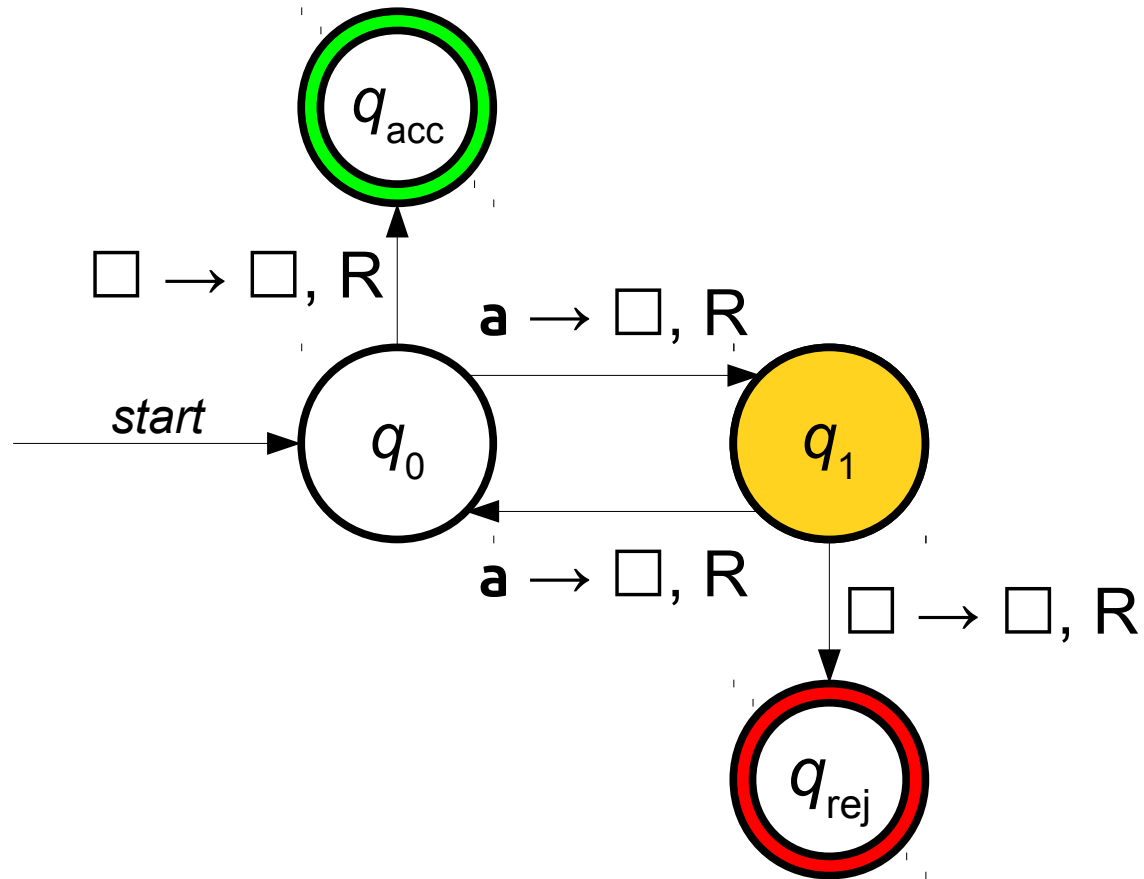
Each transition has the form

***read***  $\rightarrow$  ***write***, ***dir***

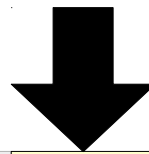
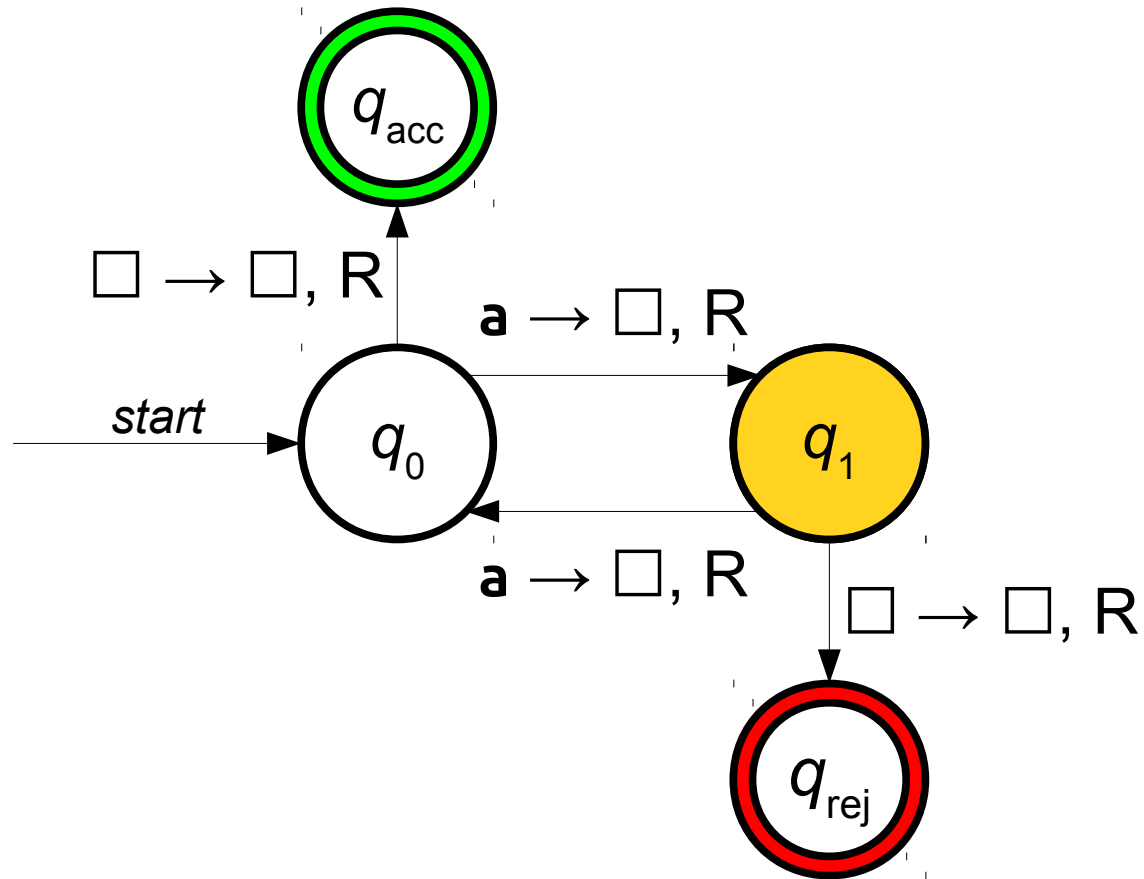
and means “if symbol ***read*** is under the tape head, replace it with ***write*** and move the tape head in direction ***dir*** (L or R). The  $\square$  symbol denotes a blank cell.



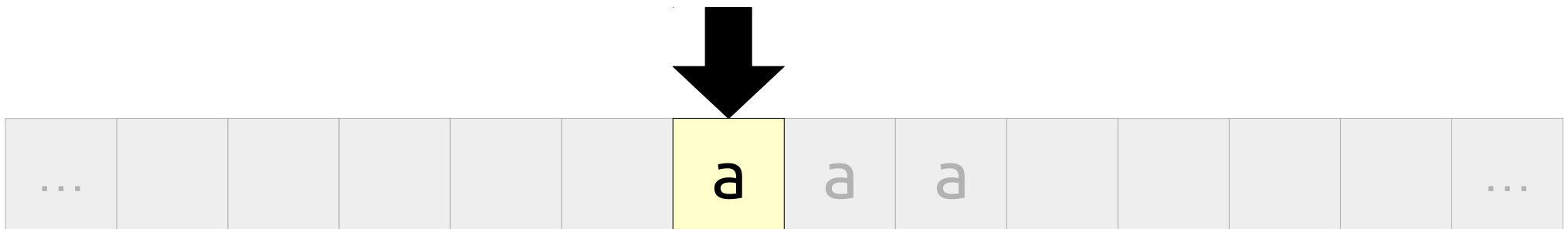
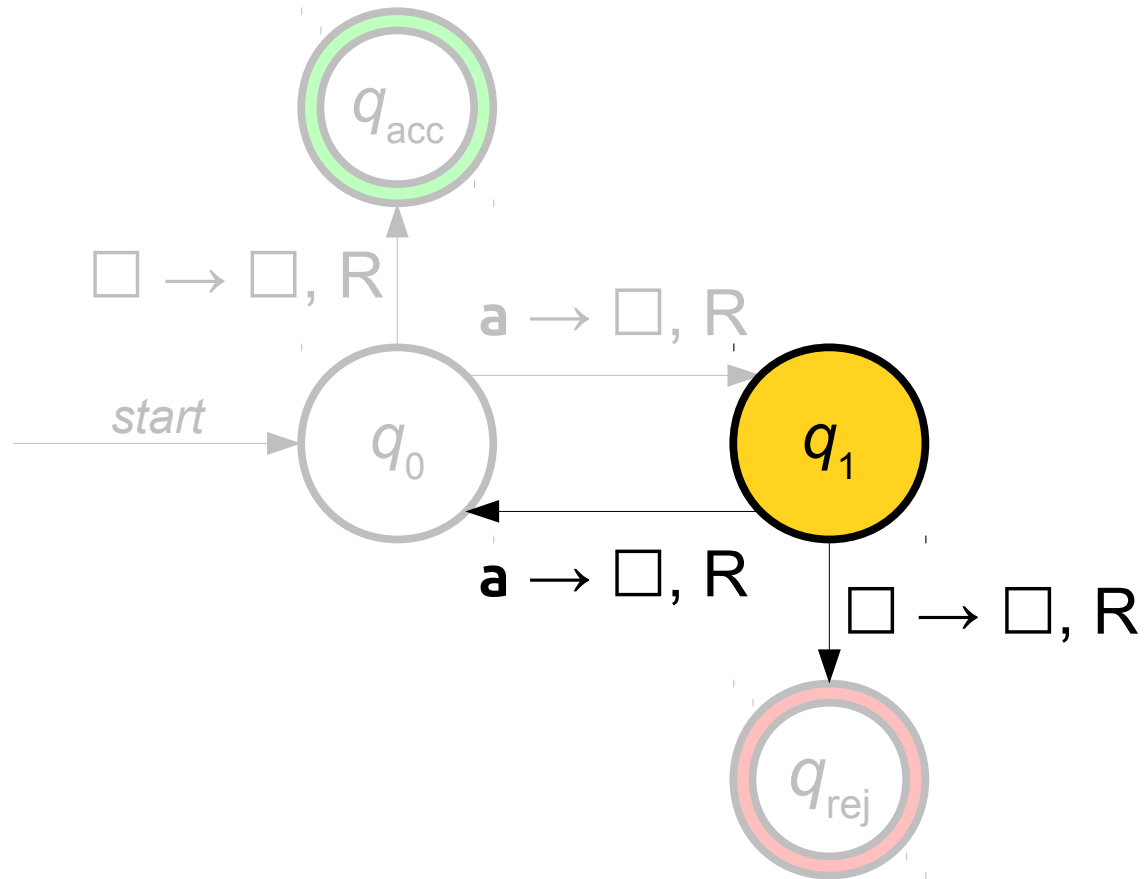
# A Simple Turing Machine



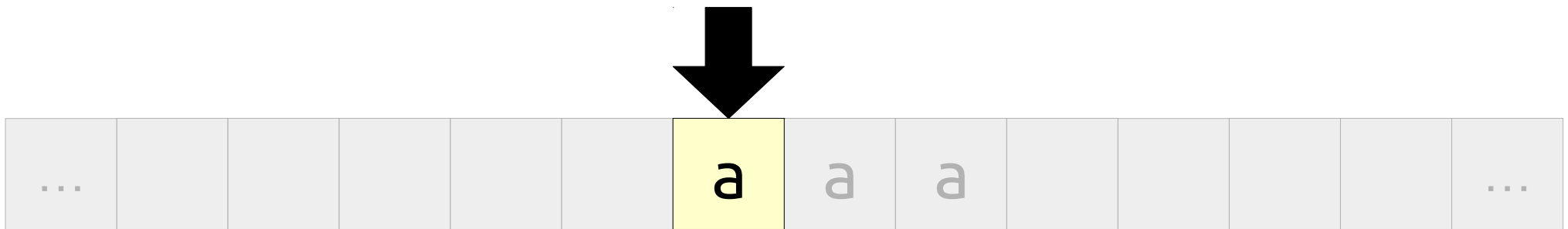
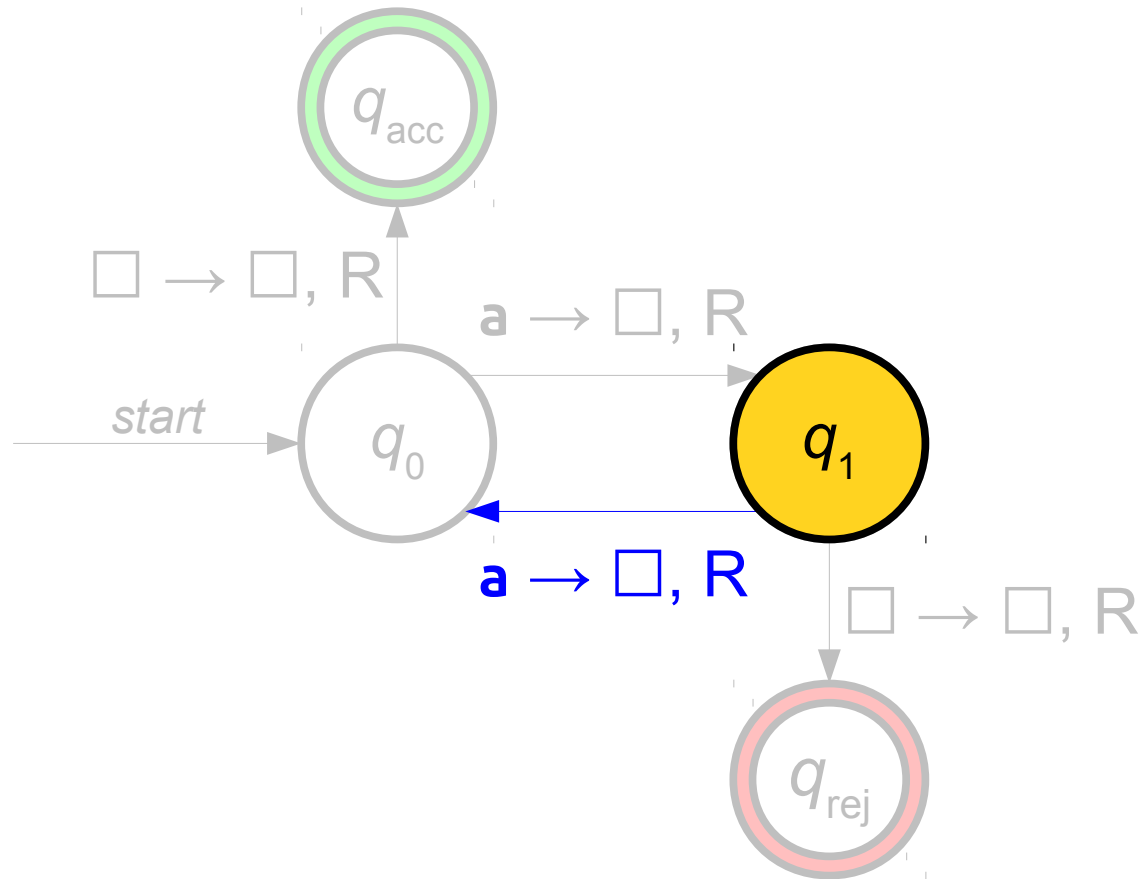
# A Simple Turing Machine



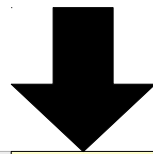
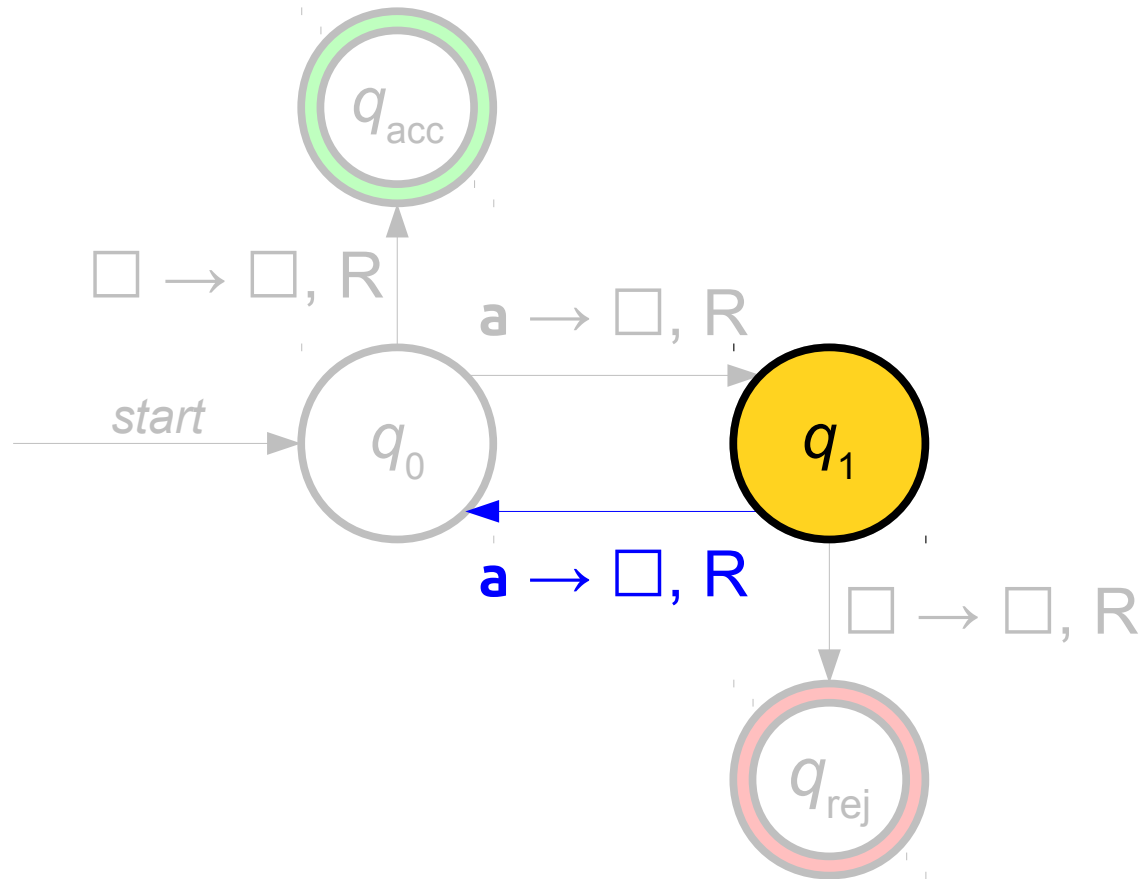
# A Simple Turing Machine



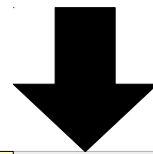
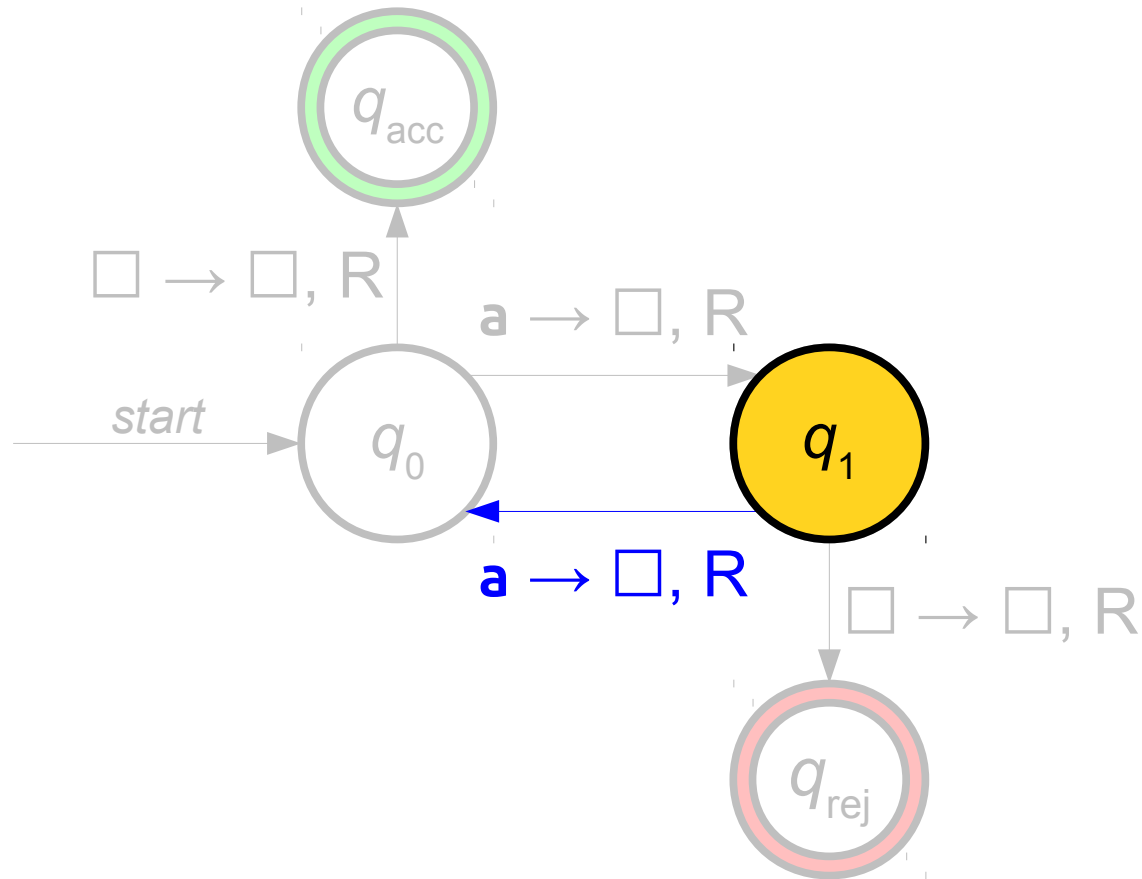
# A Simple Turing Machine



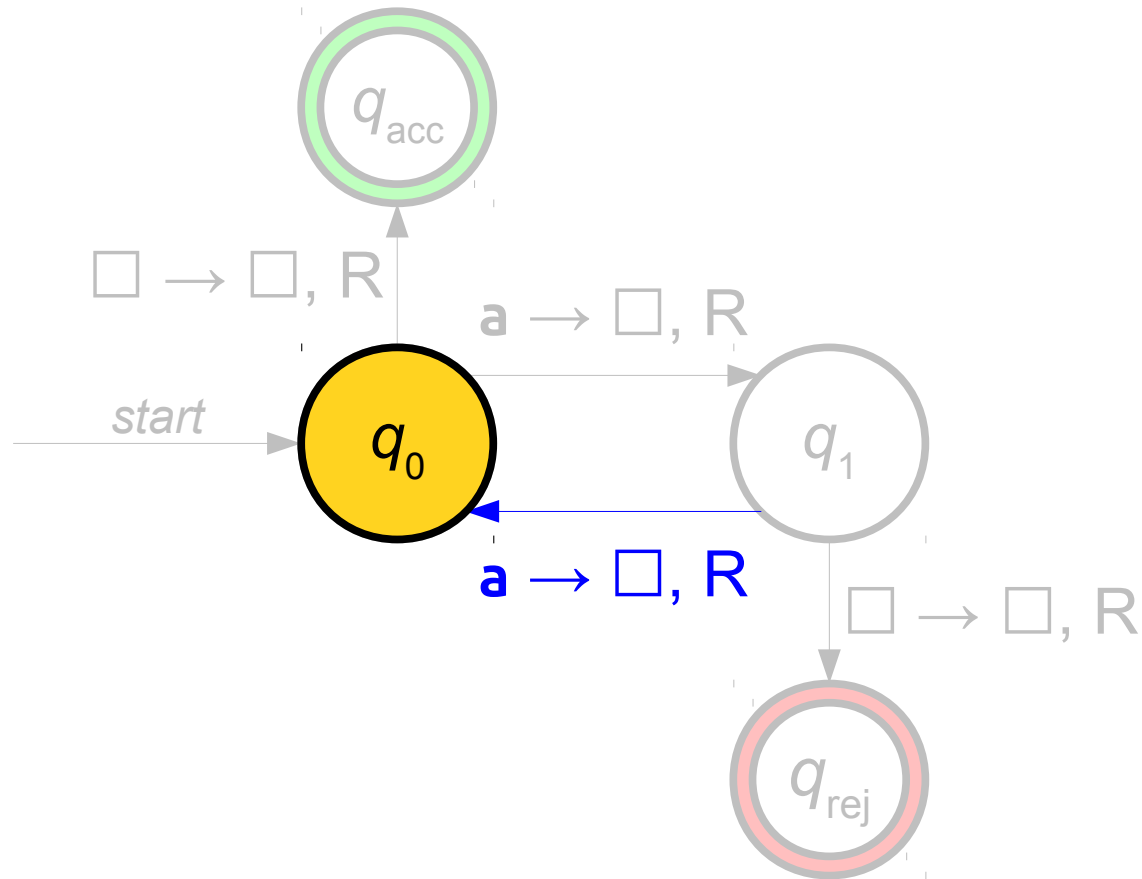
# A Simple Turing Machine



# A Simple Turing Machine

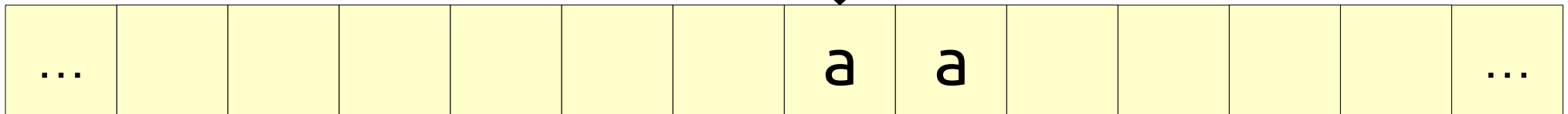
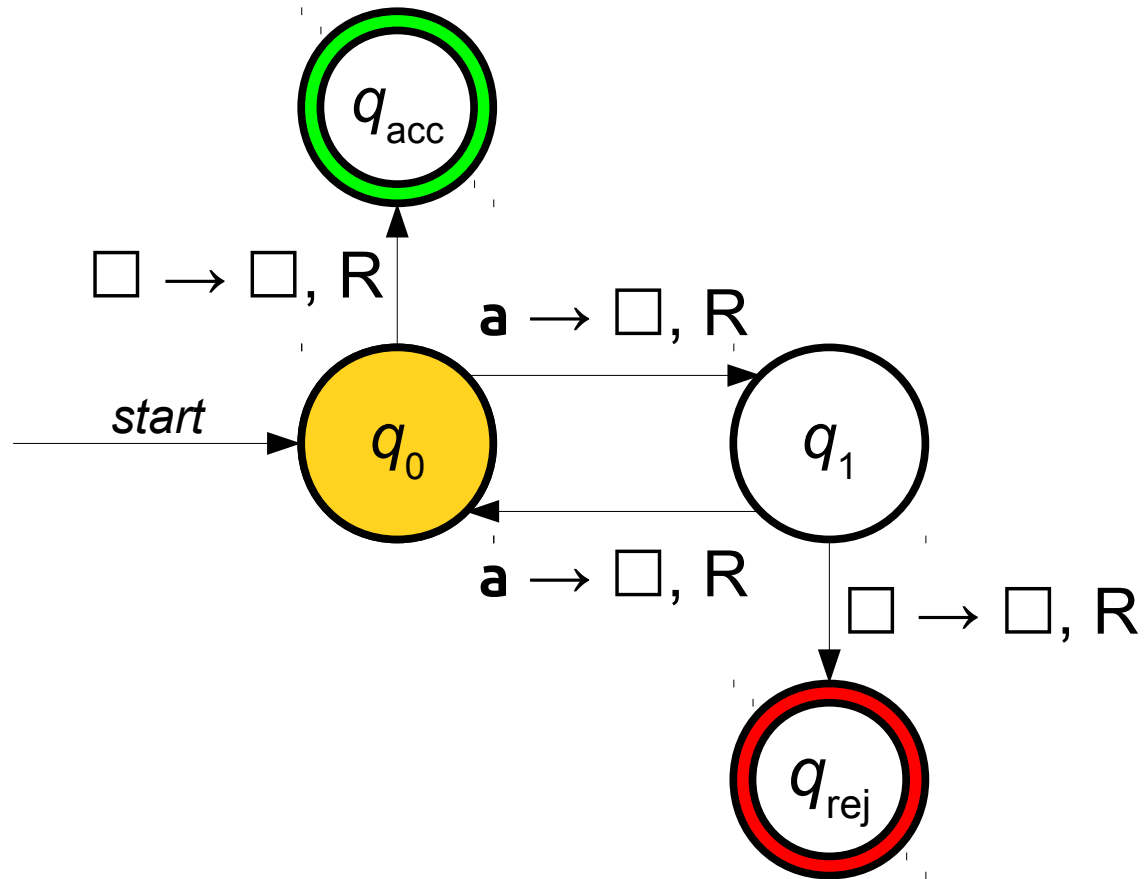


# A Simple Turing Machine

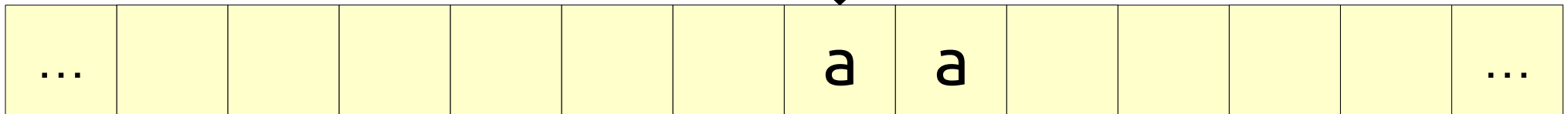
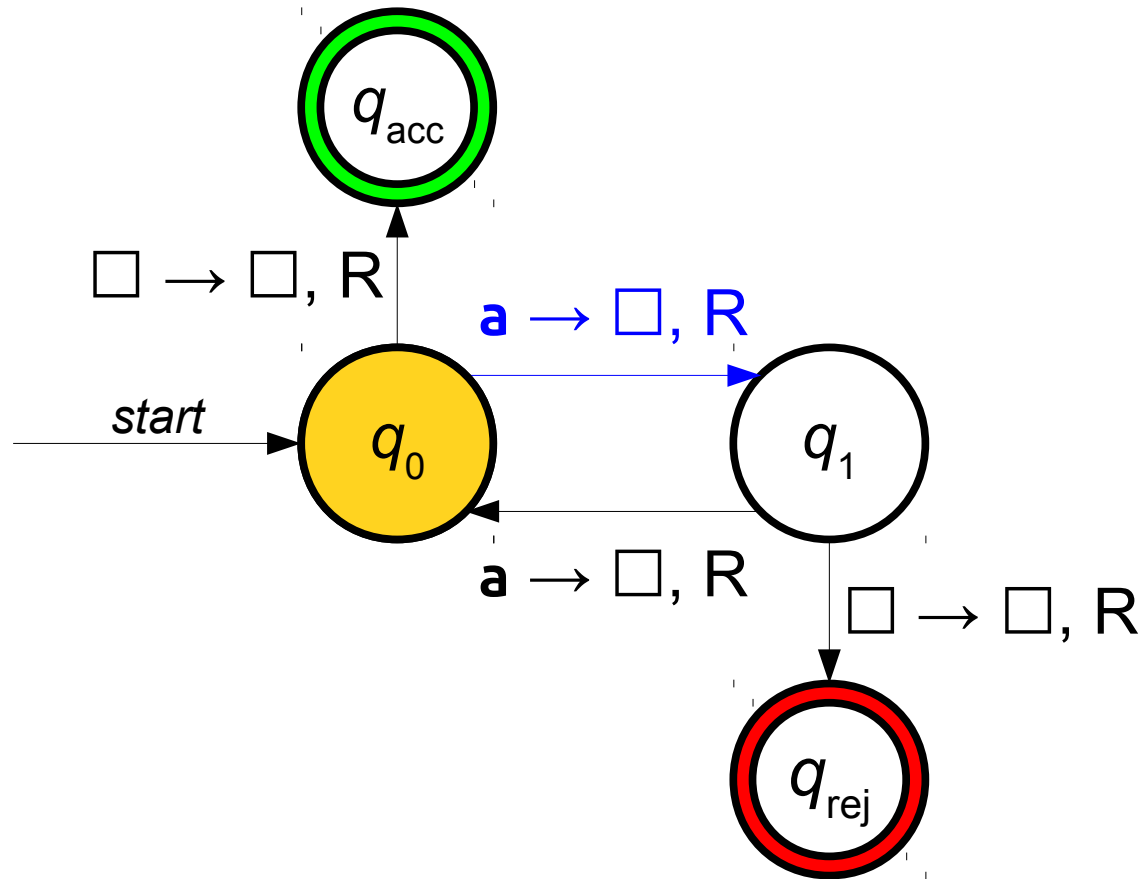




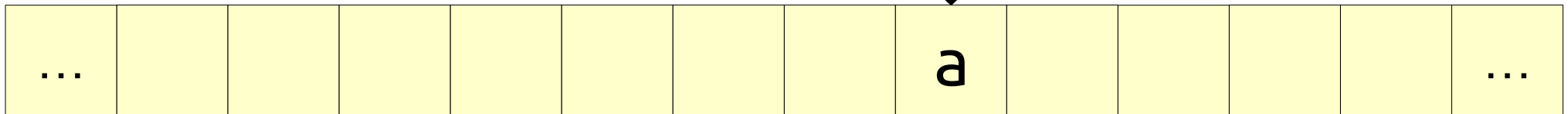
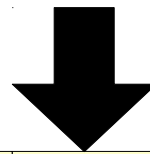
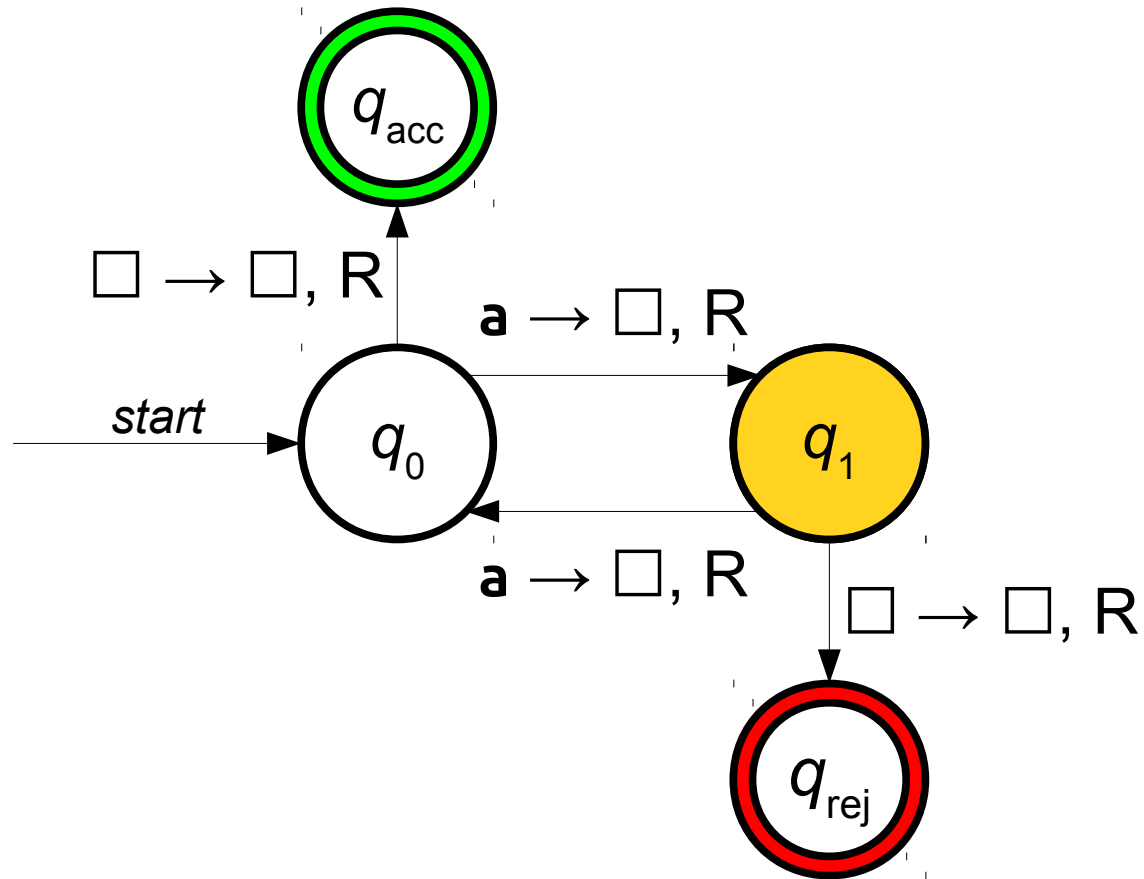
# A Simple Turing Machine



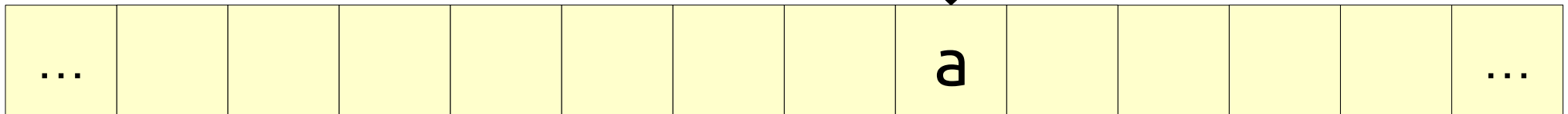
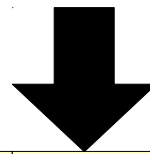
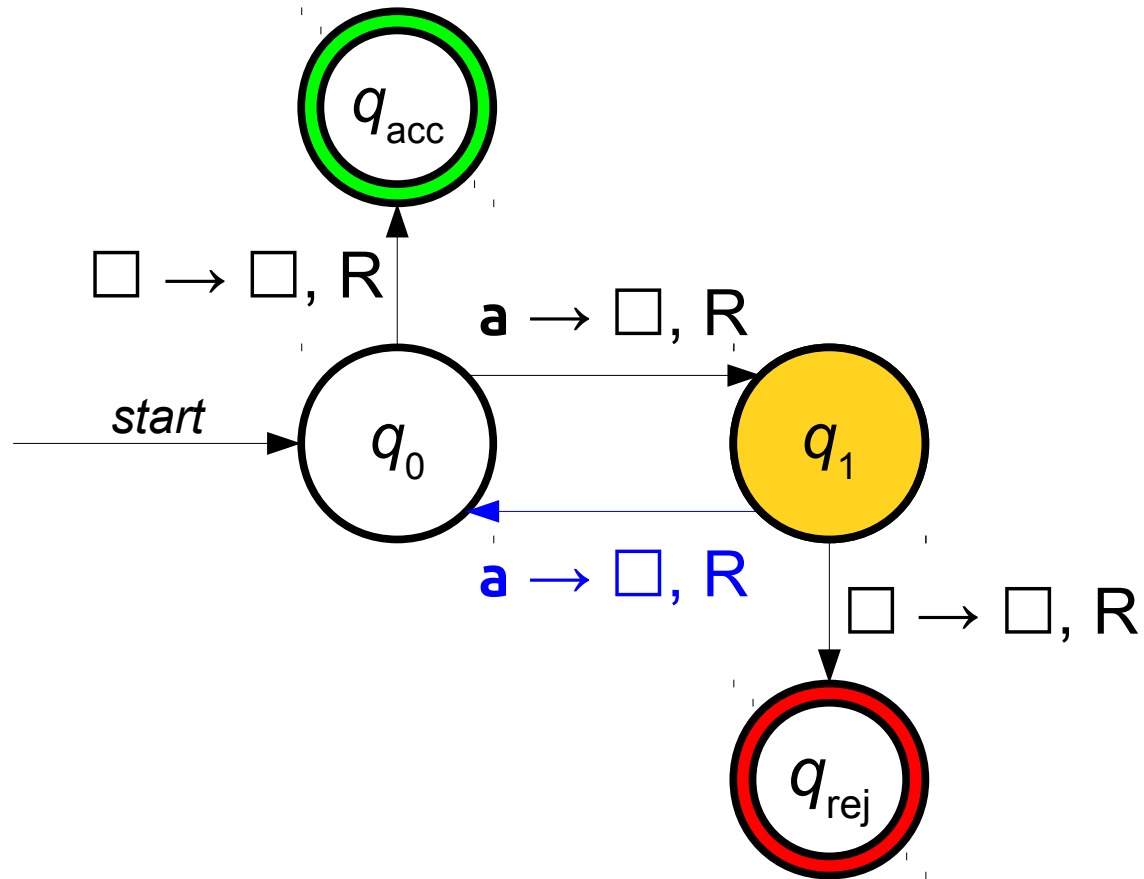
# A Simple Turing Machine



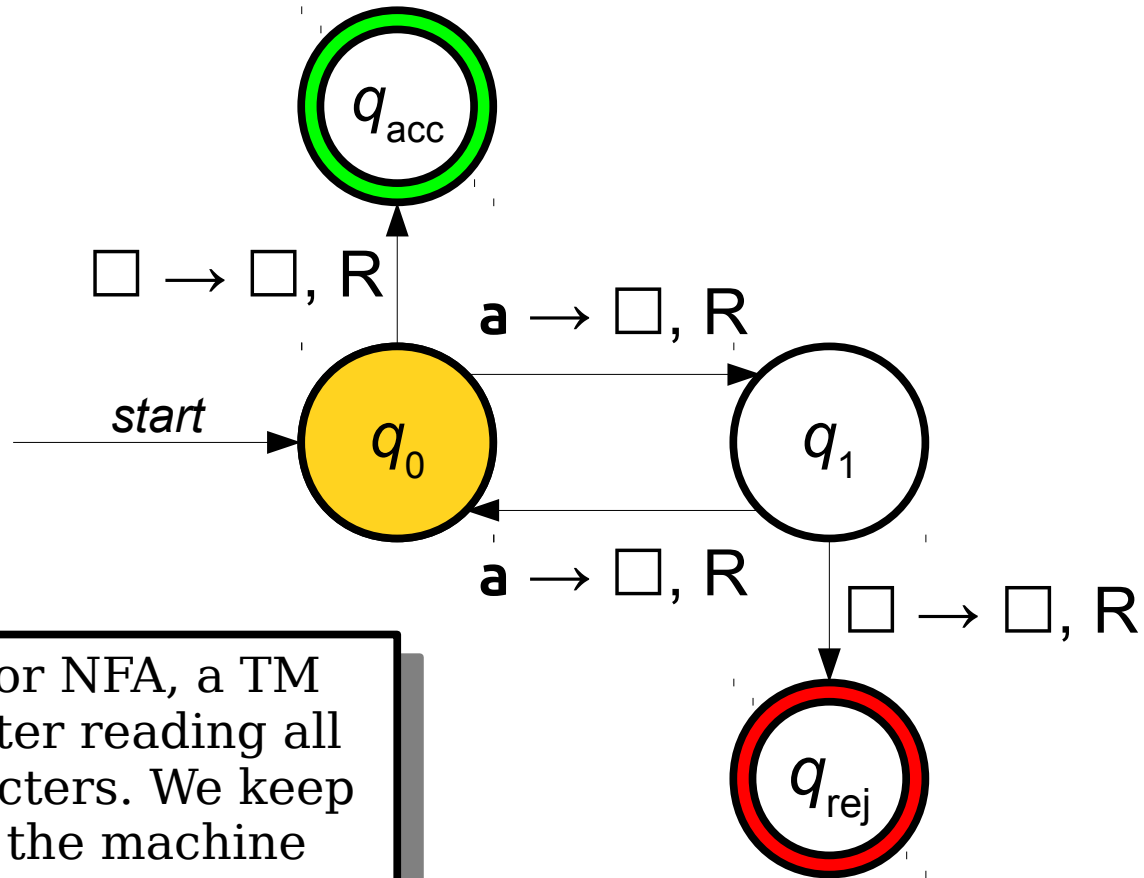
# A Simple Turing Machine



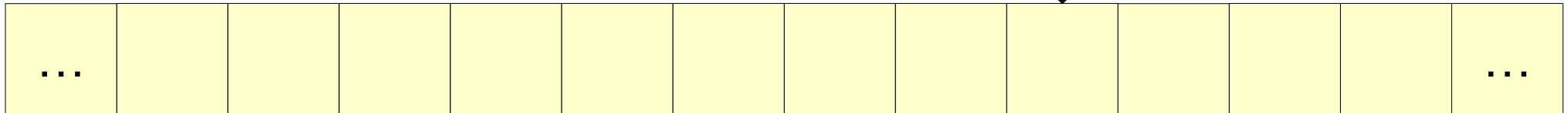
# A Simple Turing Machine



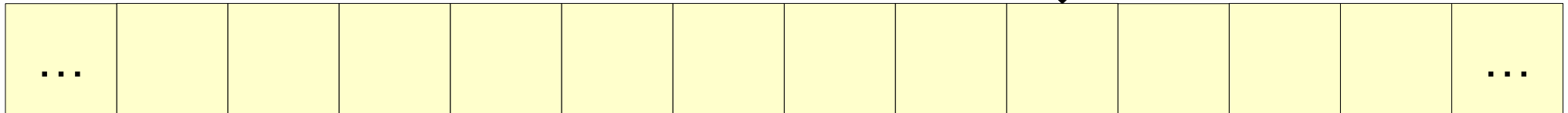
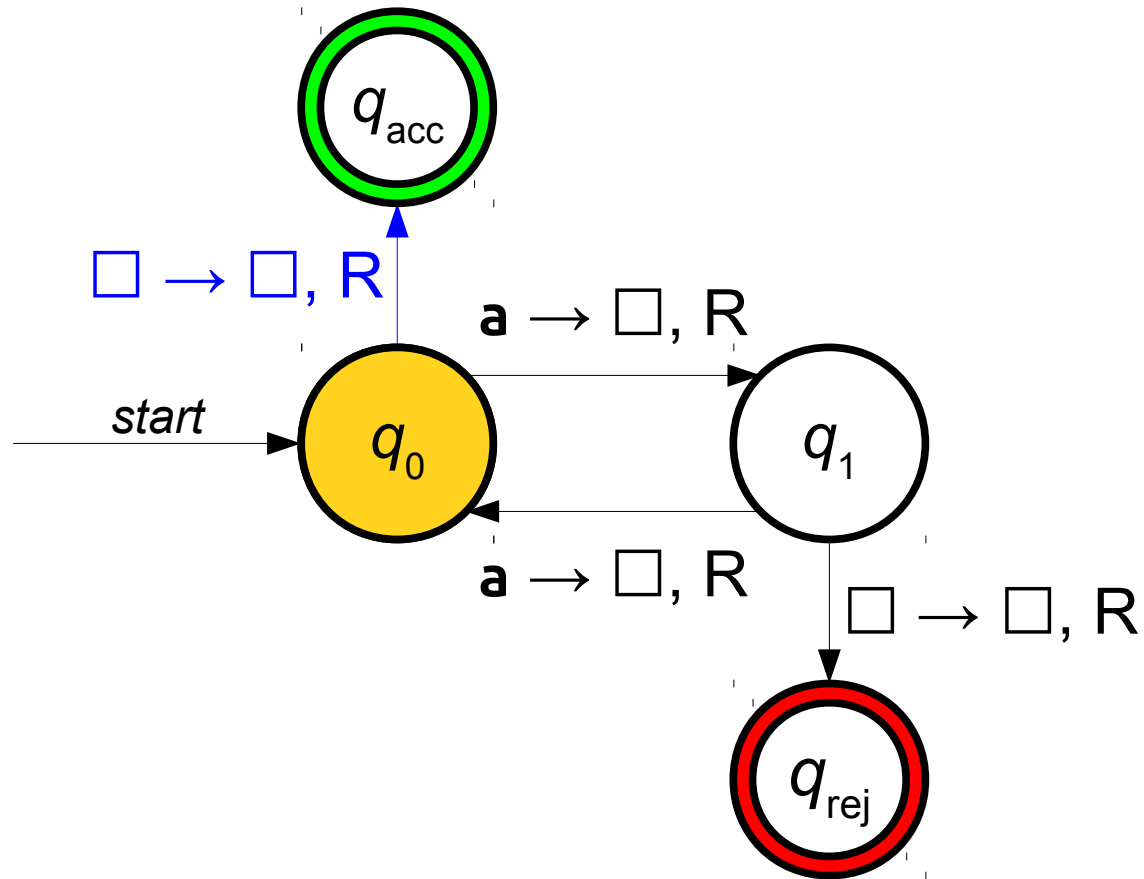
# A Simple Turing Machine



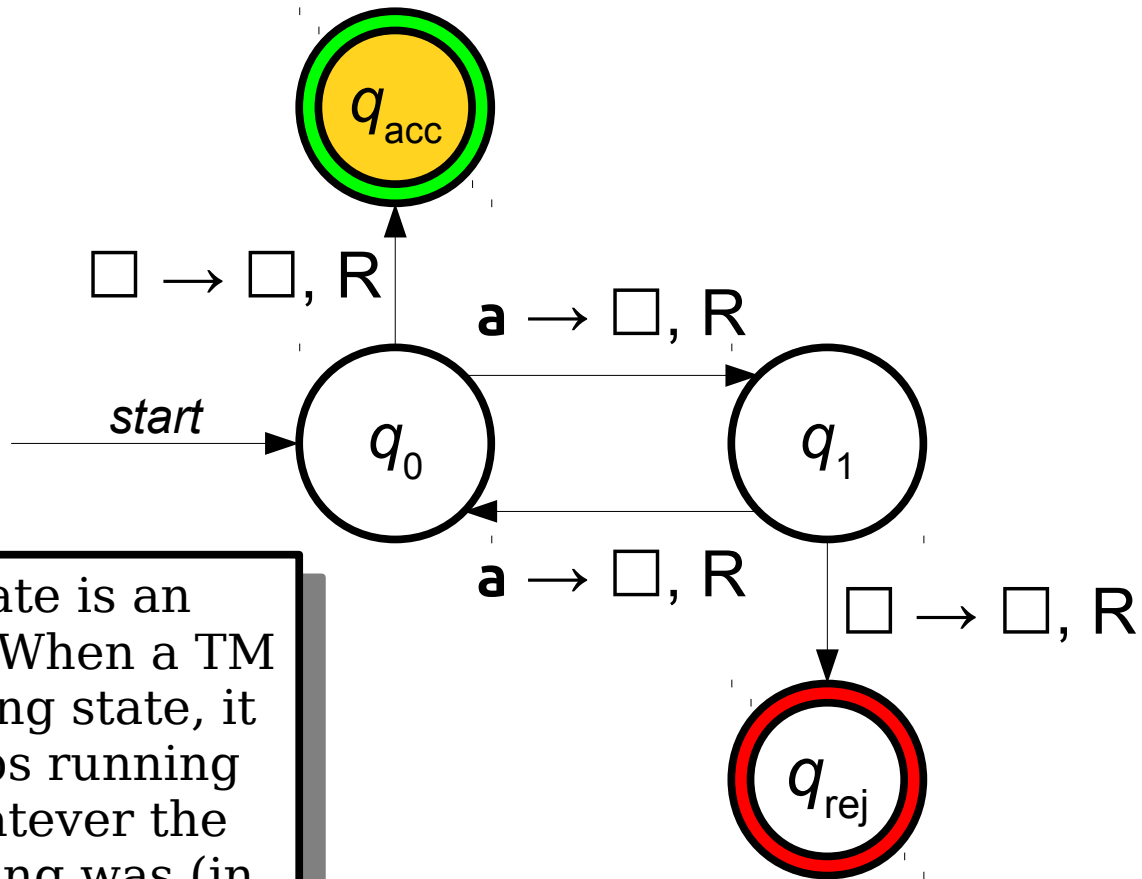
Unlike a DFA or NFA, a TM doesn't stop after reading all the input characters. We keep running until the machine explicitly says to stop.



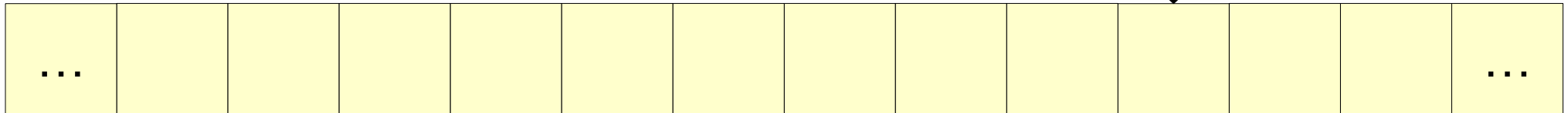
# A Simple Turing Machine



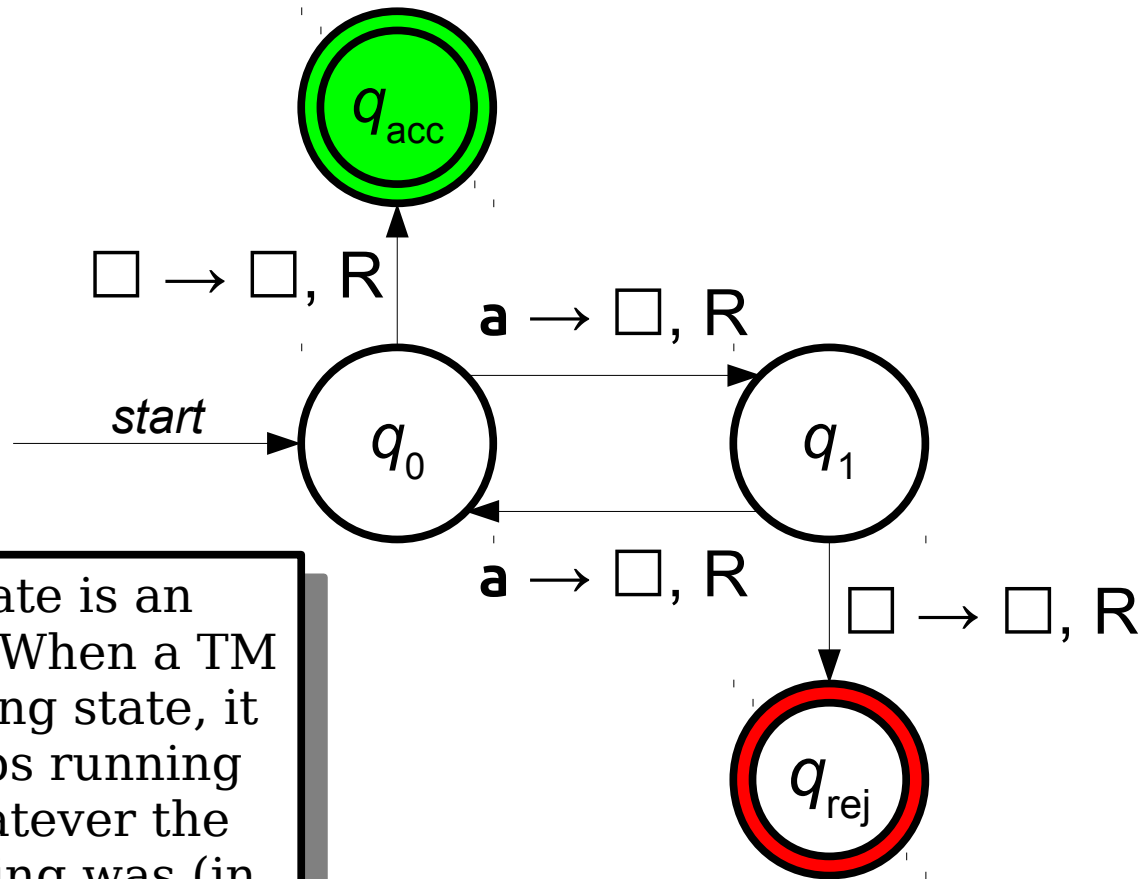
# A Simple Turing Machine



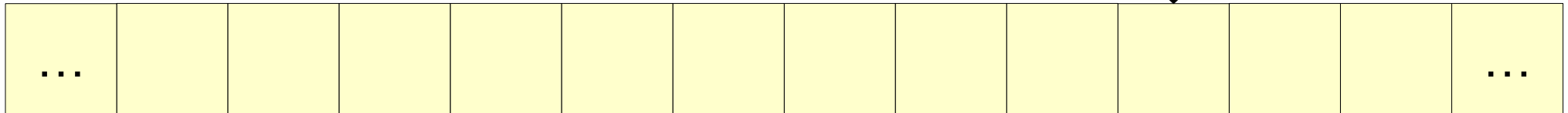
This special state is an **accepting state**. When a TM enters an accepting state, it *immediately* stops running and accepts whatever the original input string was (in this case, **aaaa**).



# A Simple Turing Machine

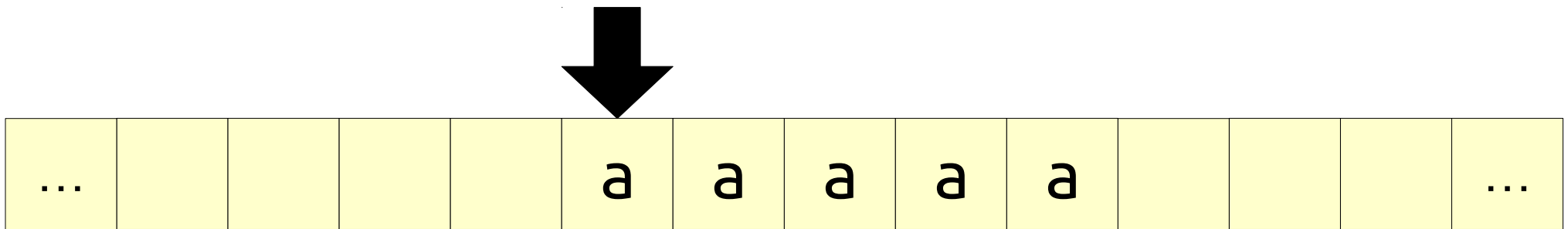
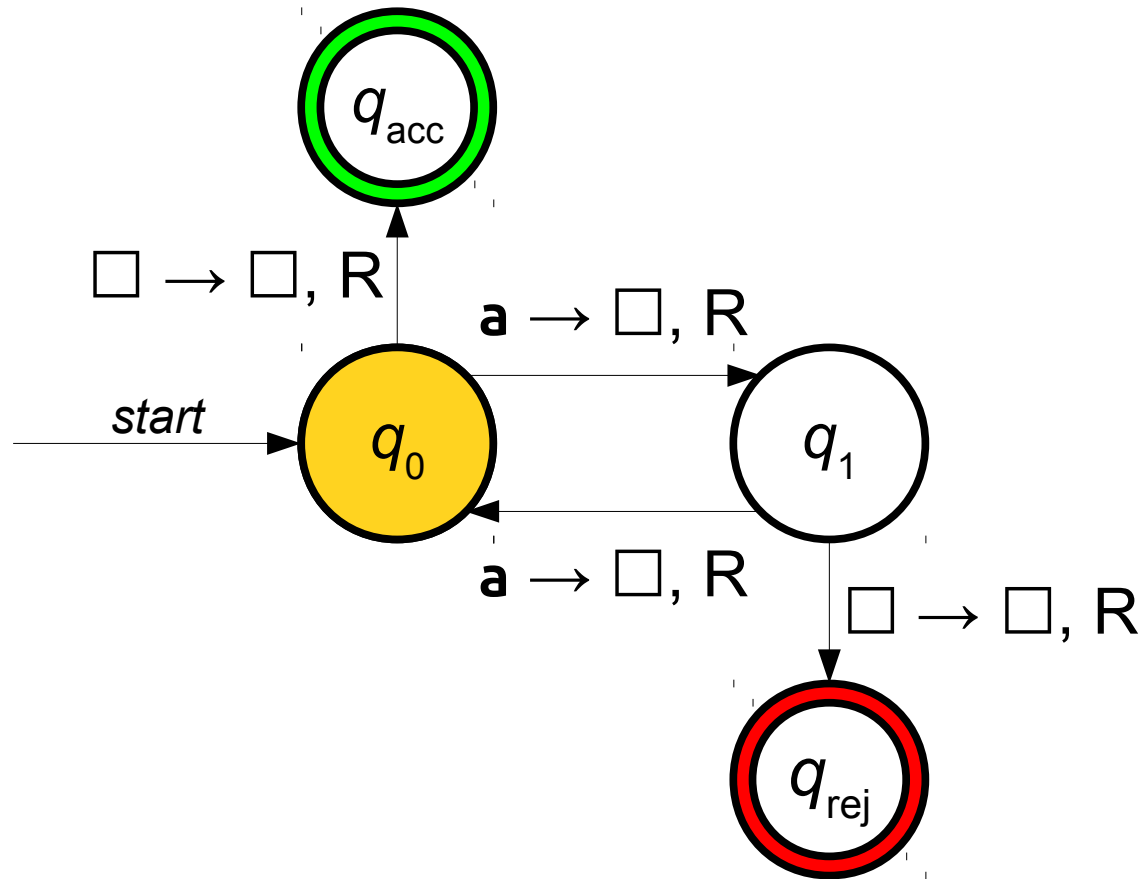


This special state is an **accepting state**. When a TM enters an accepting state, it *immediately* stops running and accepts whatever the original input string was (in this case, **aaaa**).

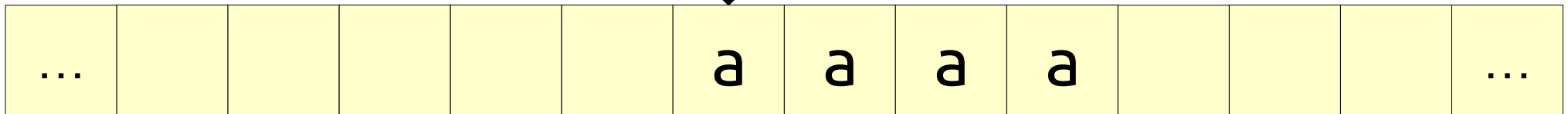
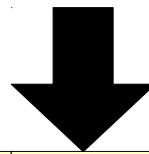
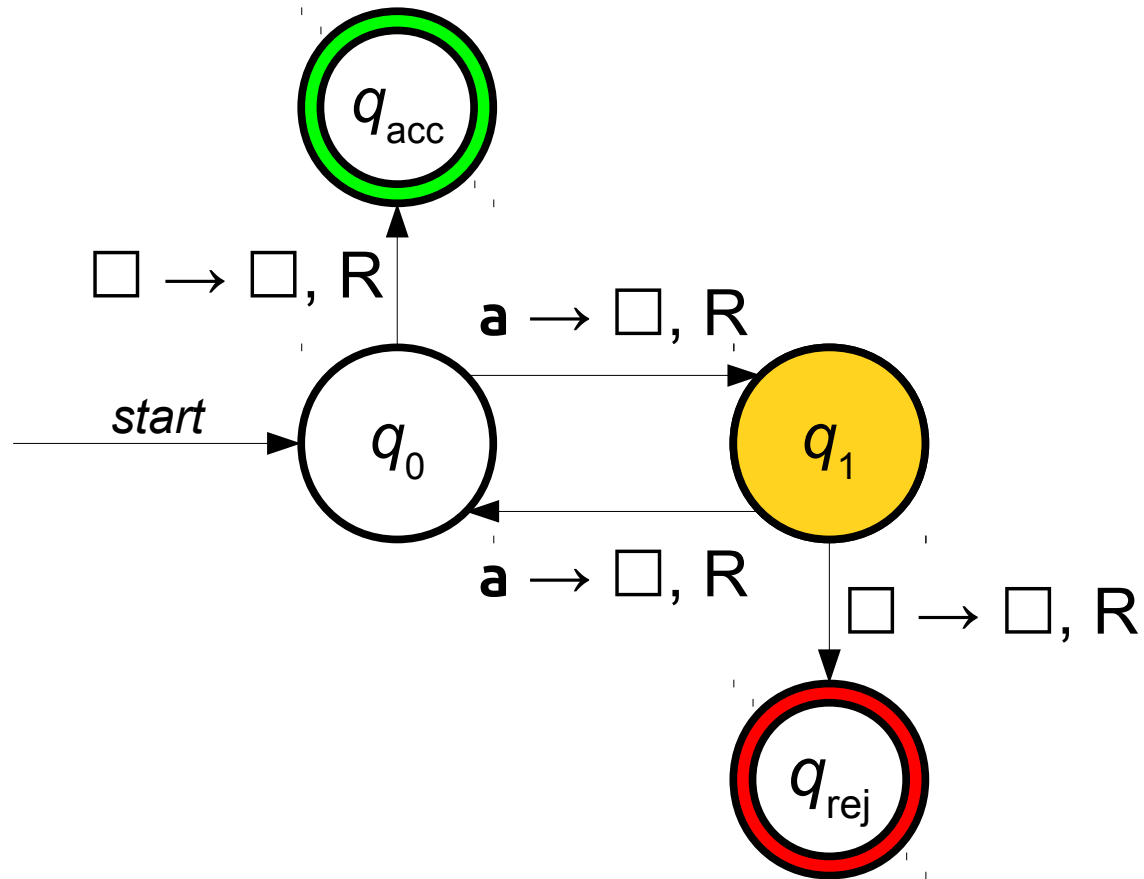




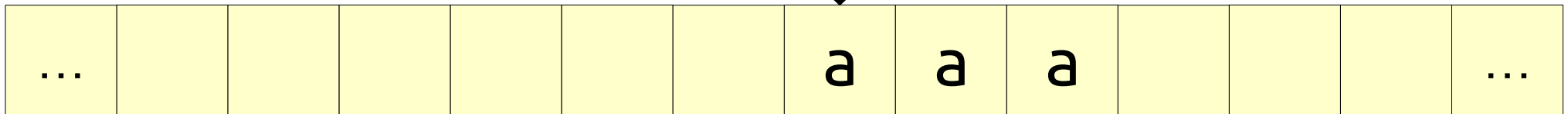
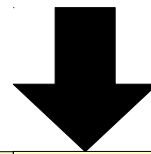
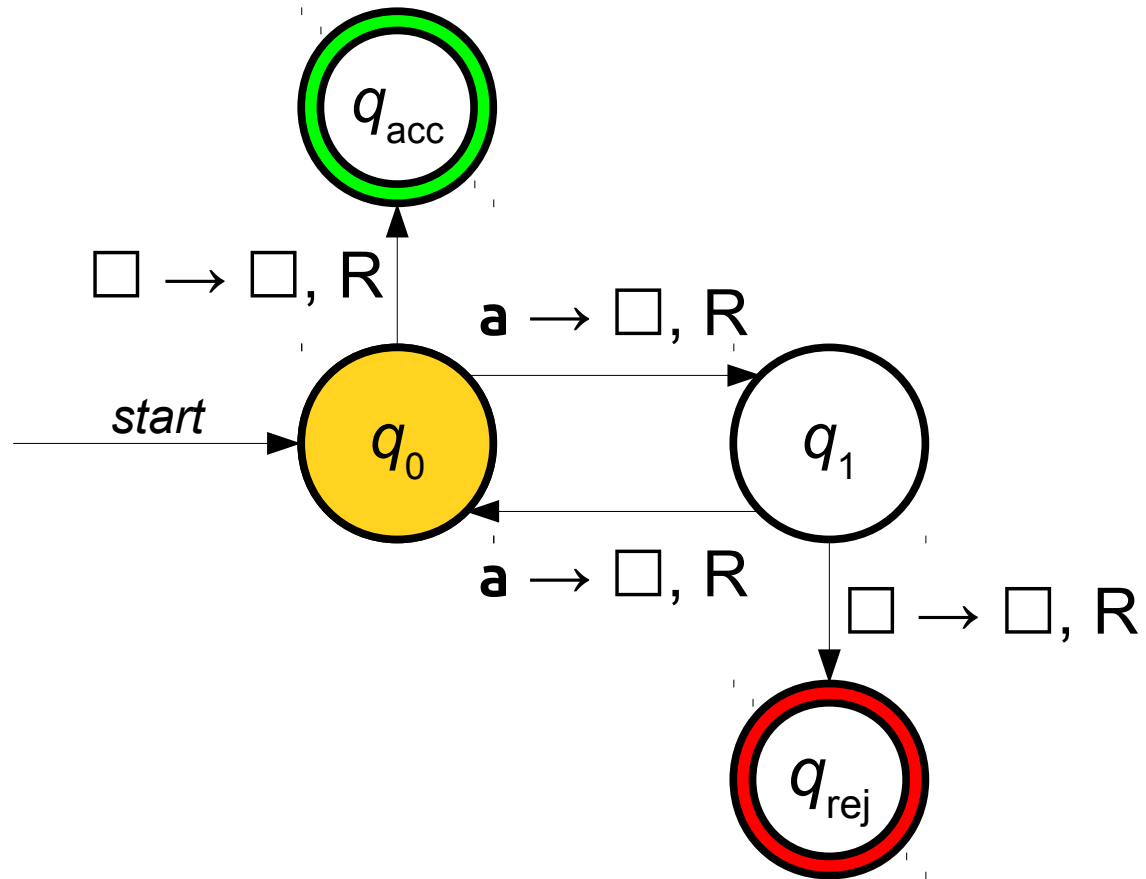
# A Simple Turing Machine



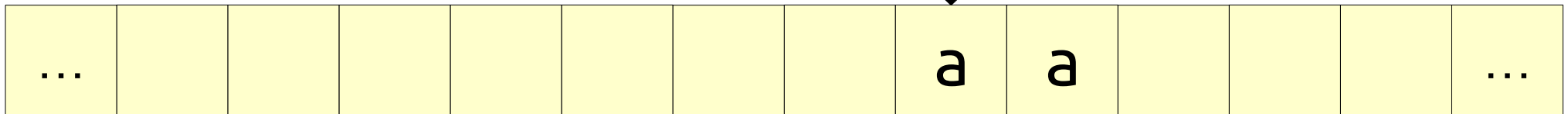
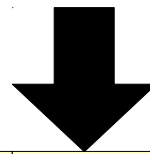
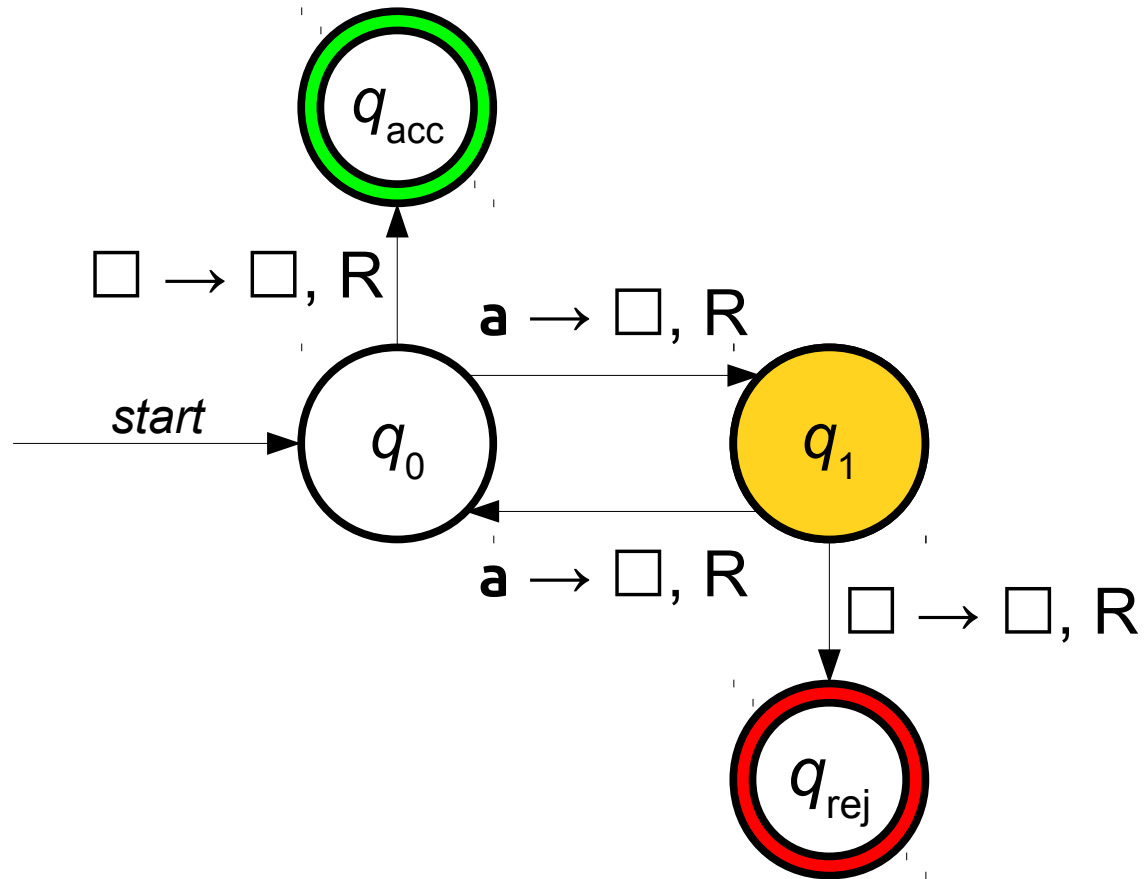
# A Simple Turing Machine



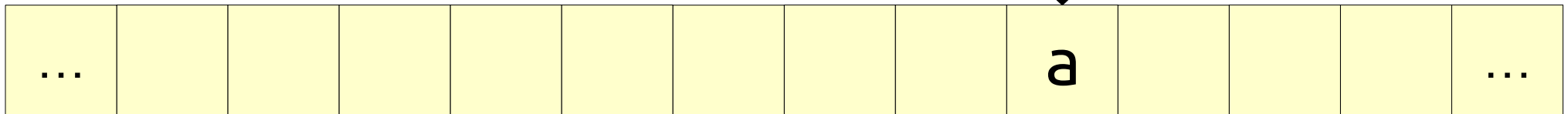
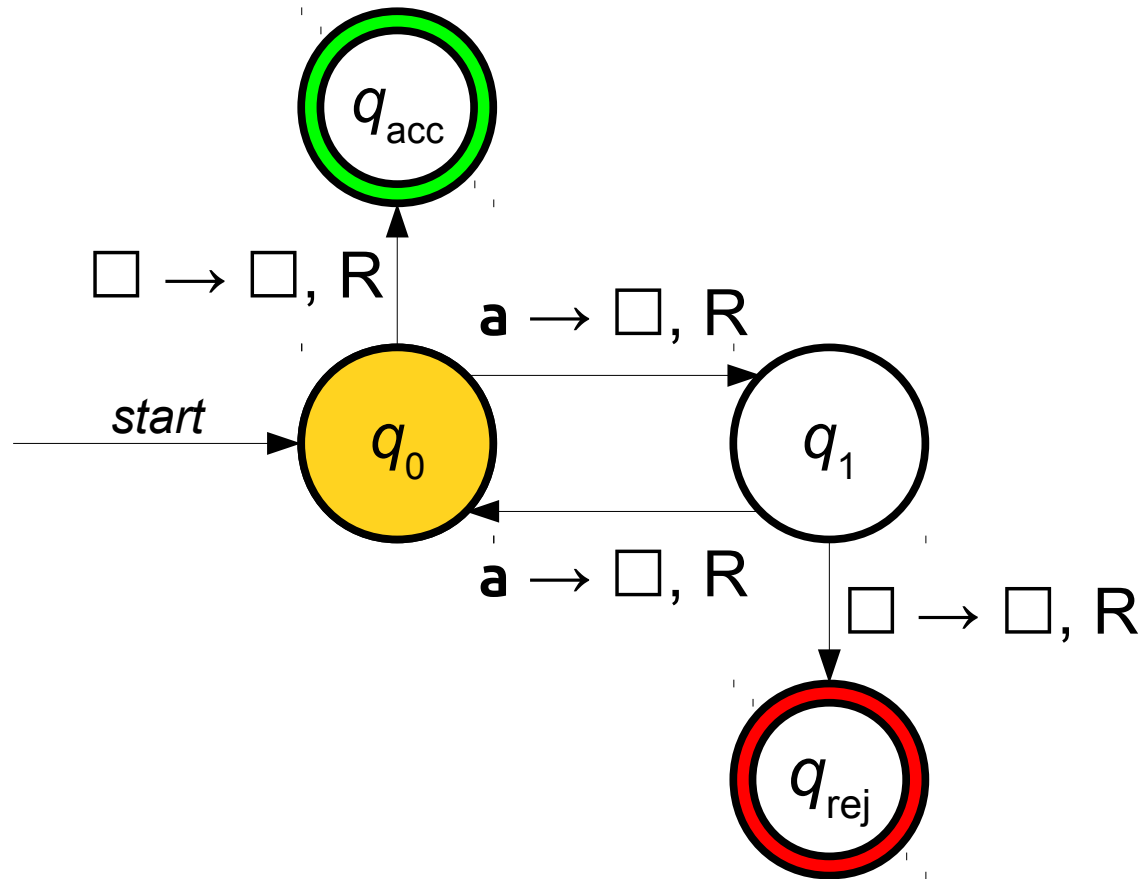
# A Simple Turing Machine



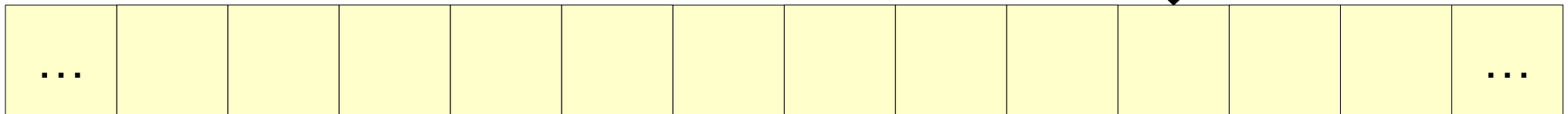
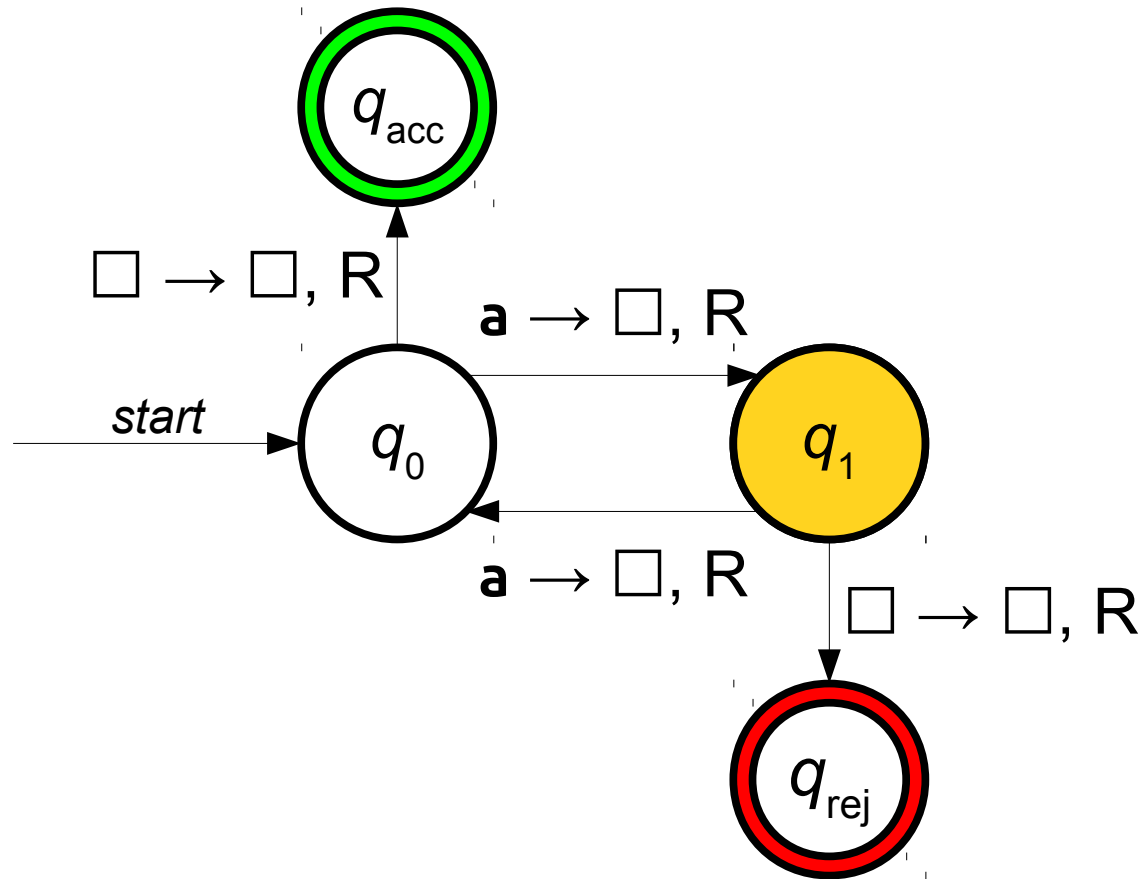
# A Simple Turing Machine



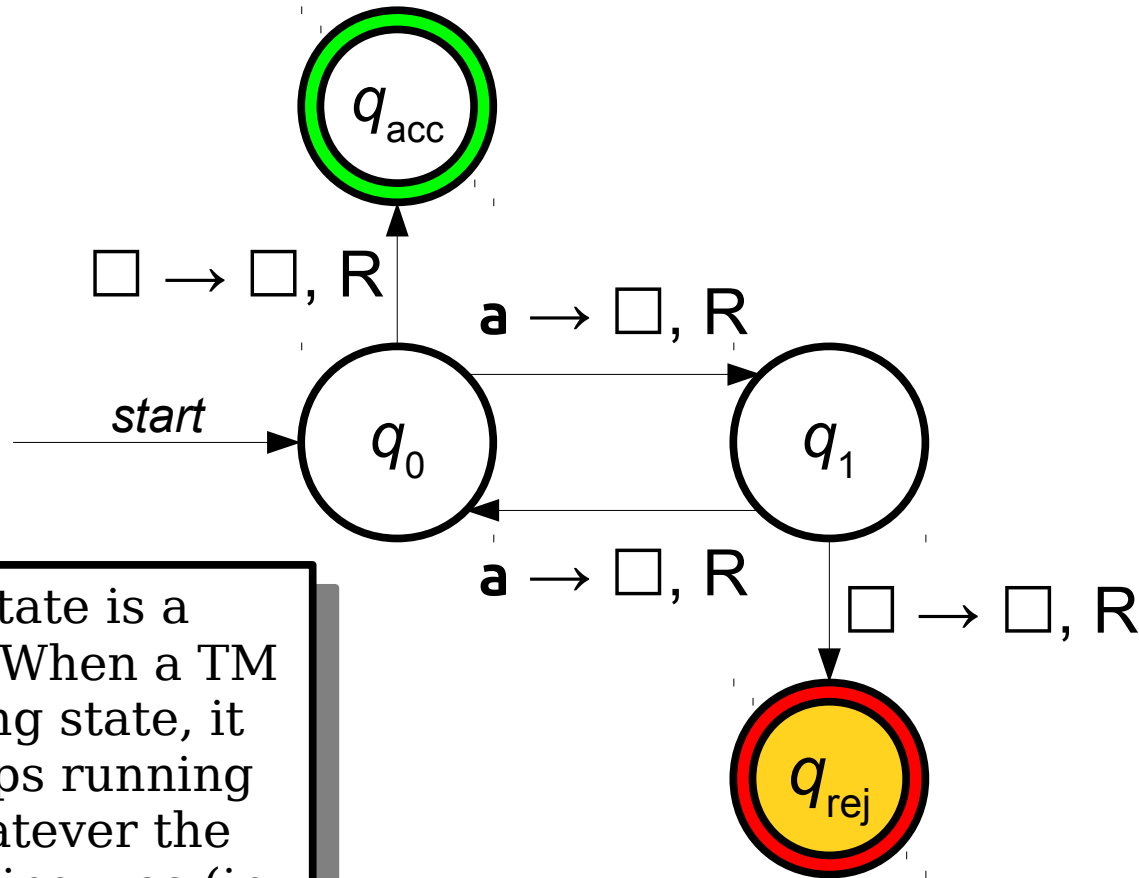
# A Simple Turing Machine



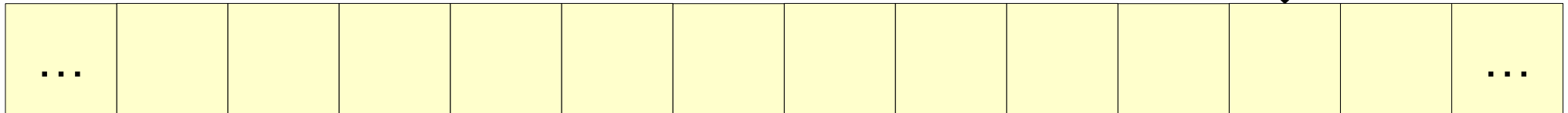
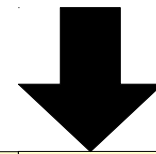
# A Simple Turing Machine



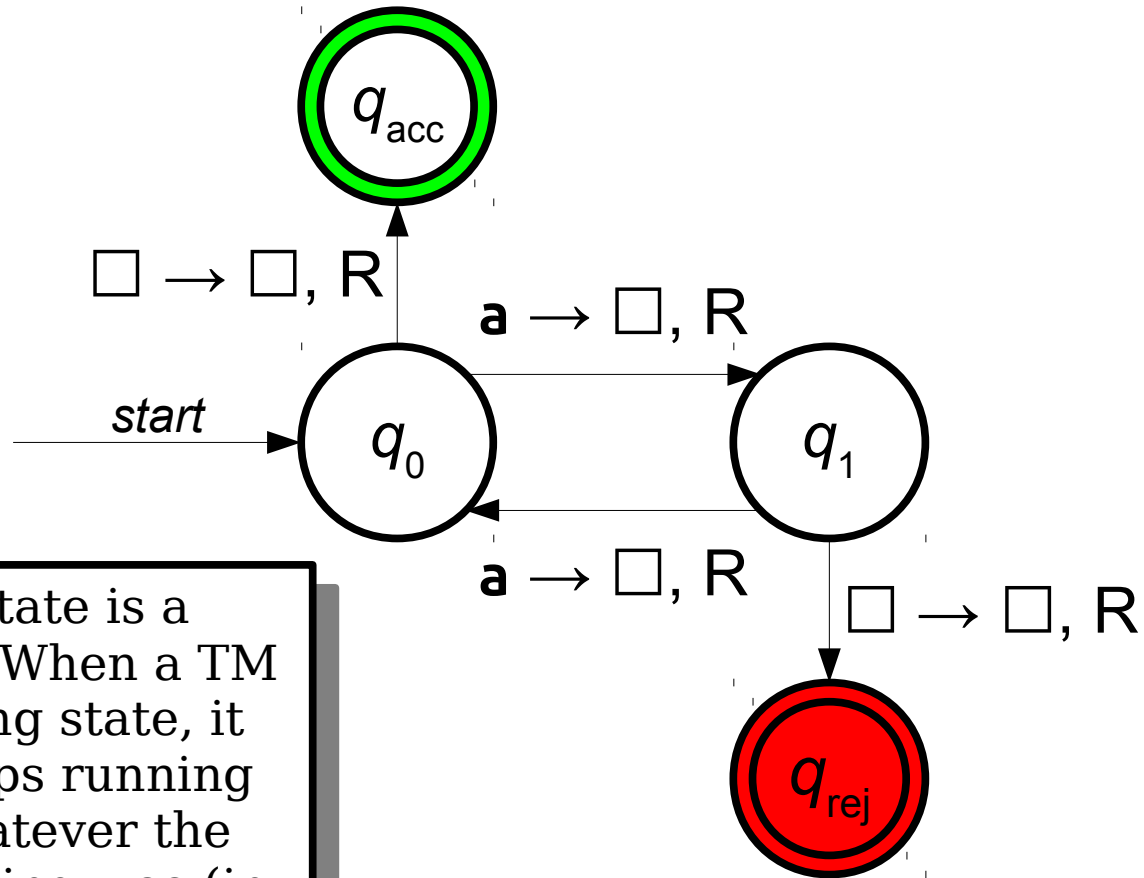
# A Simple Turing Machine



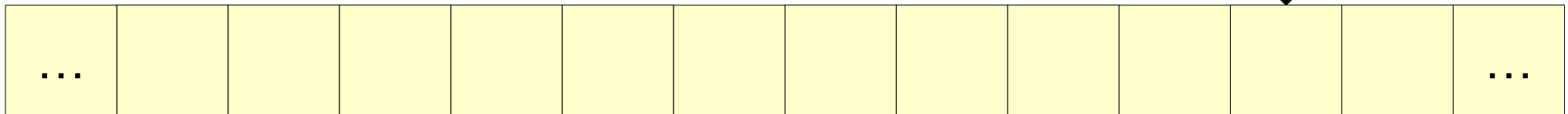
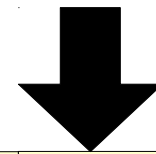
This special state is a **rejecting state**. When a TM enters a rejecting state, it *immediately* stops running and rejects whatever the original input string was (in this case, **aaaaa**).



# A Simple Turing Machine

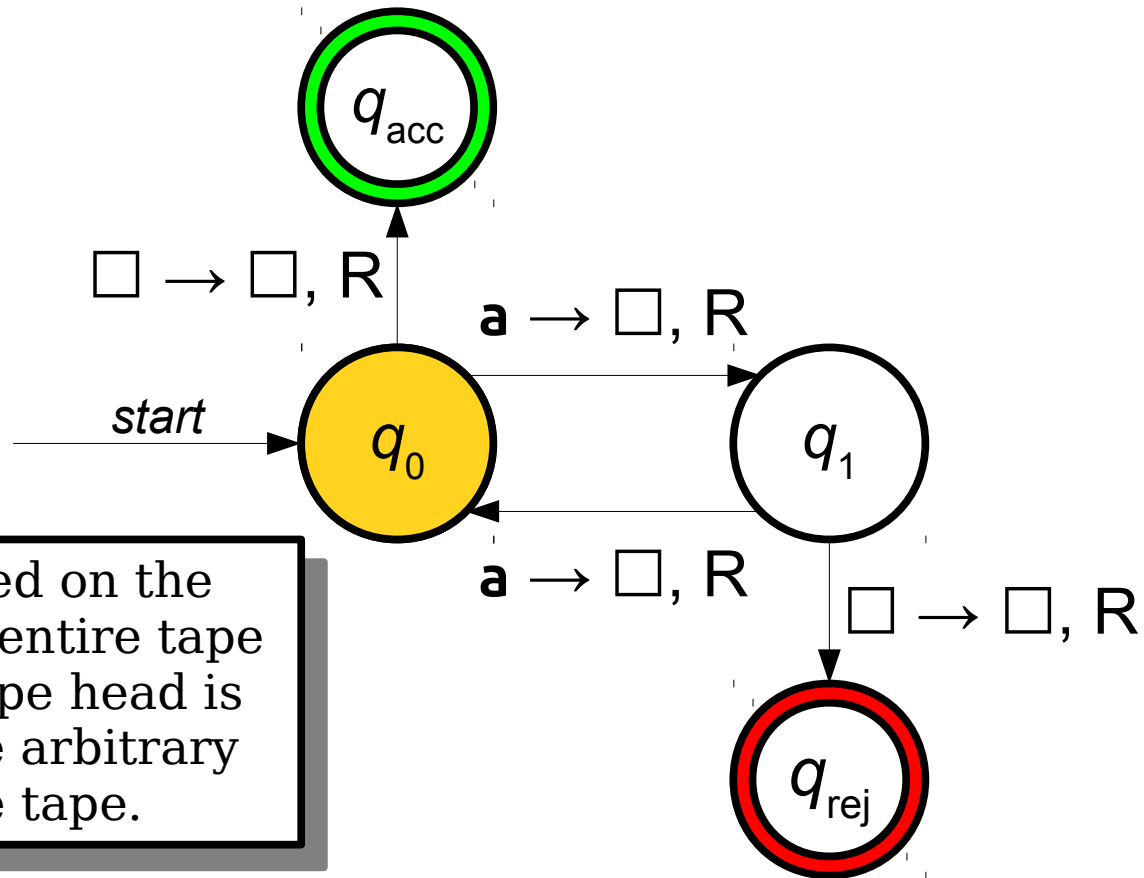


This special state is a **rejecting state**. When a TM enters a rejecting state, it *immediately* stops running and rejects whatever the original input string was (in this case, **aaaaa**).

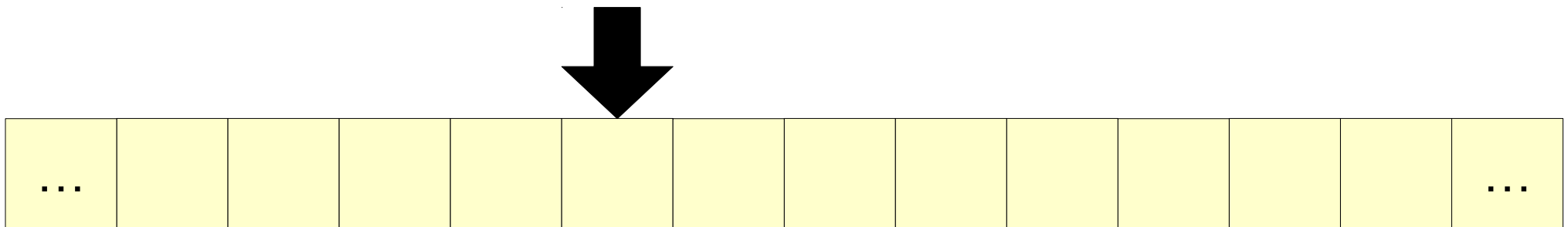




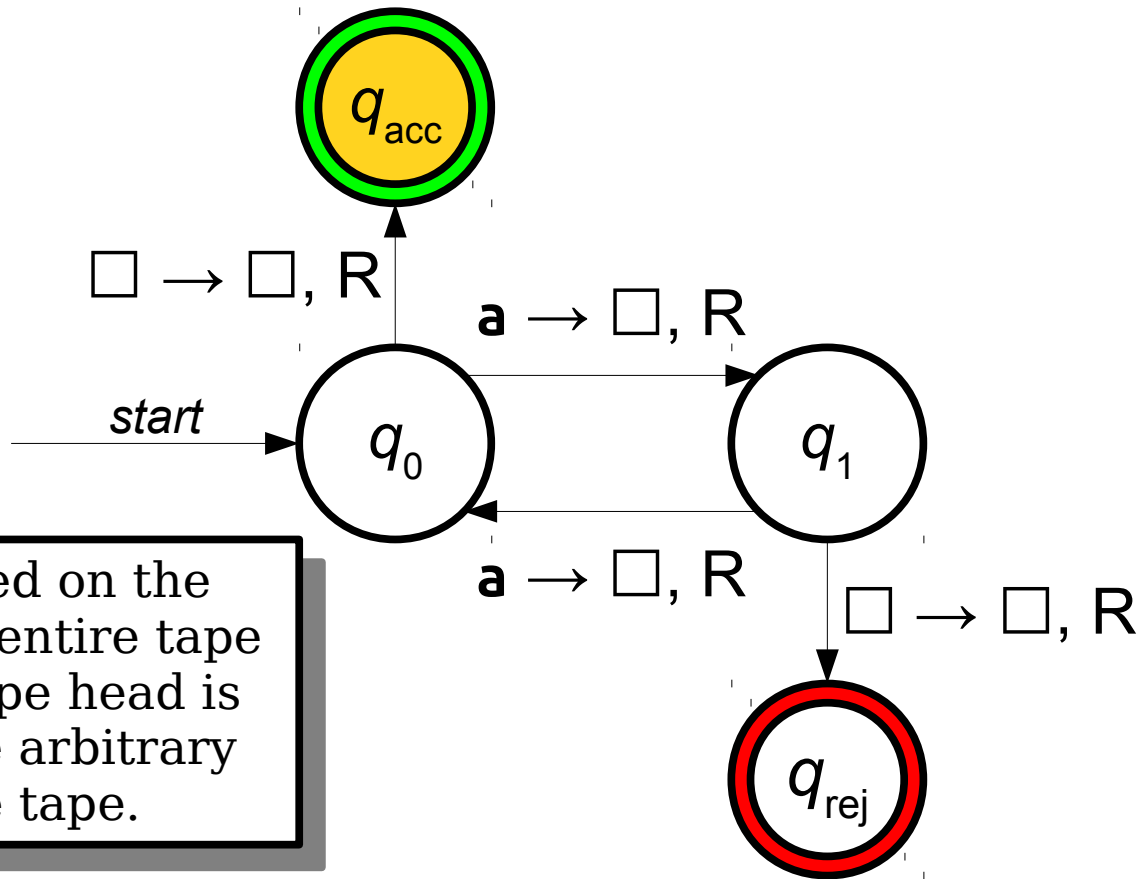
# A Simple Turing Machine



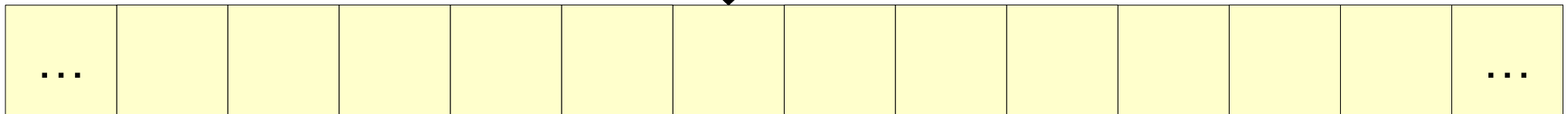
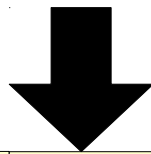
If the TM is started on the empty string  $\varepsilon$ , the entire tape is blank and the tape head is positioned at some arbitrary location on the tape.



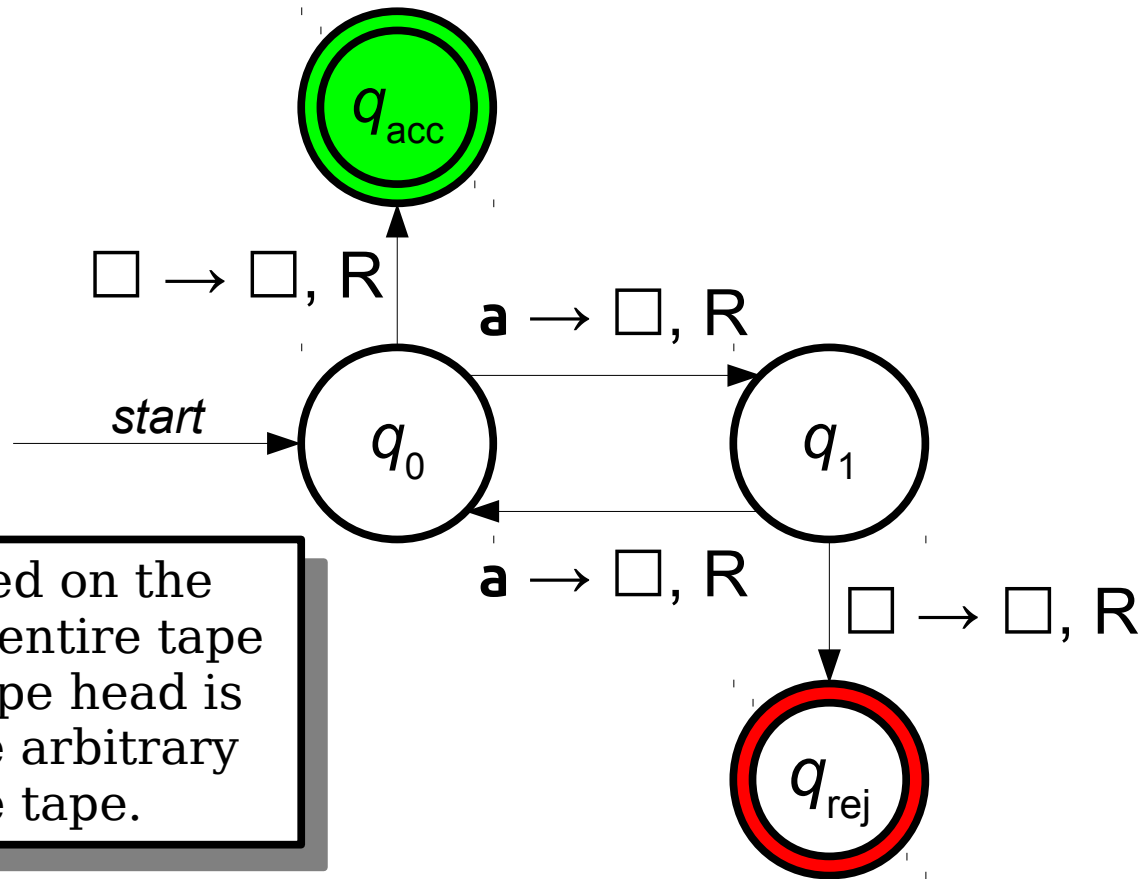
# A Simple Turing Machine



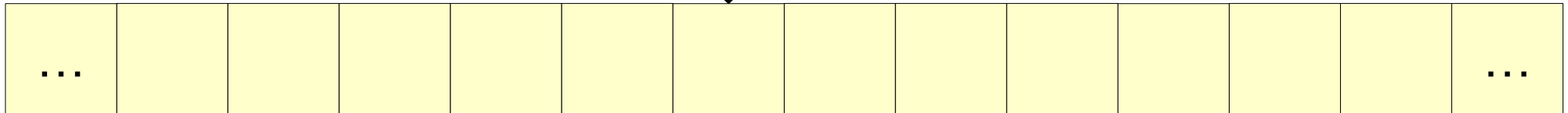
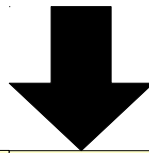
If the TM is started on the empty string  $\varepsilon$ , the entire tape is blank and the tape head is positioned at some arbitrary location on the tape.



# A Simple Turing Machine



If the TM is started on the empty string  $\varepsilon$ , the entire tape is blank and the tape head is positioned at some arbitrary location on the tape.



# The Turing Machine

- A Turing machine consists of three parts:
  - A ***finite-state control*** that issues commands,
  - an ***infinite tape*** for input and scratch space, and
  - a ***tape head*** that can read and write a single tape cell.
- At each step, the Turing machine
  - writes a symbol to the tape cell under the tape head,
  - changes state, and
  - moves the tape head to the left or to the right.

# Input and Tape Alphabets

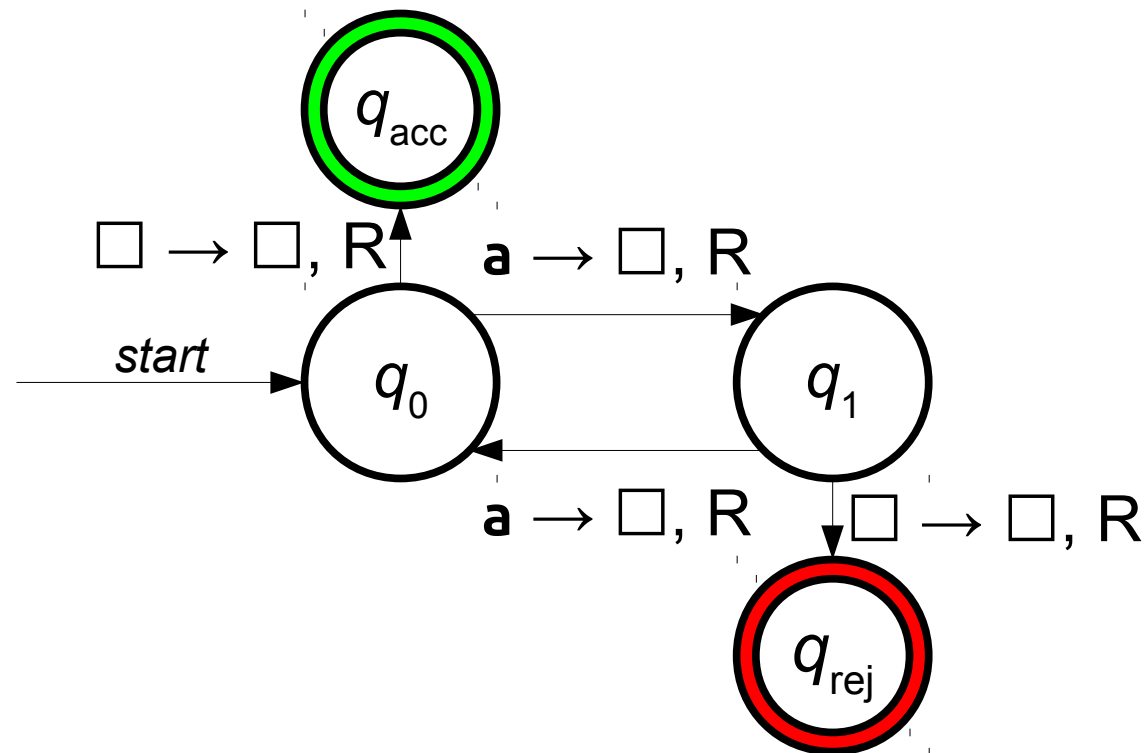
- A Turing machine has two alphabets:
  - An **input alphabet**  $\Sigma$ . All input strings are written in the input alphabet.
  - A **tape alphabet**  $\Gamma$ , where  $\Sigma \subsetneq \Gamma$ . The tape alphabet contains all symbols that can be written onto the tape.
- The tape alphabet  $\Gamma$  can contain any number of symbols, but always contains at least one **blank symbol**, denoted  $\square$ . You are guaranteed  $\square \notin \Sigma$ .
- At startup, the Turing machine begins with an infinite tape of  $\square$  symbols with the input written at some location. The tape head is positioned at the start of the input.

# Accepting and Rejecting States

- Unlike DFAs, Turing machines do not stop processing the input when they finish reading it.
- Turing machines decide when (and if!) they will accept or reject their input.
- Turing machines can enter infinite loops and never accept or reject; more on that later...

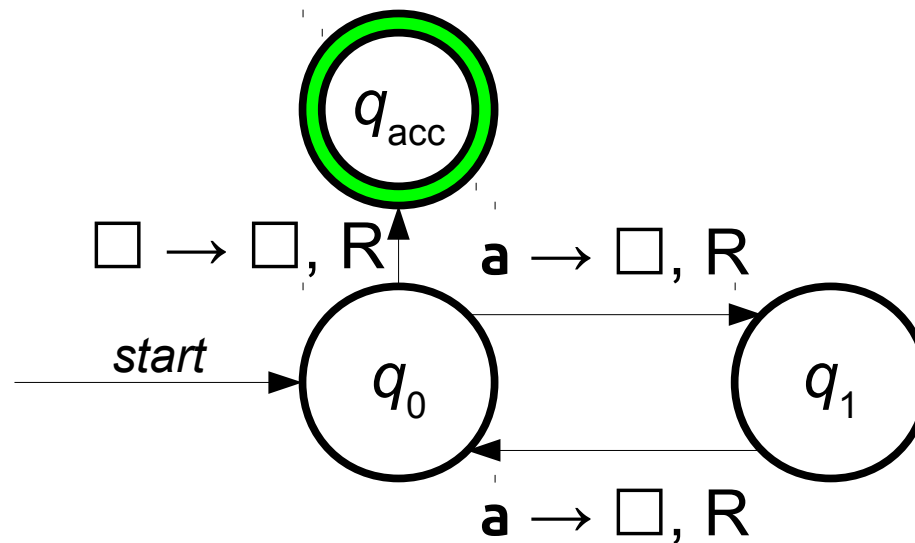
# Determinism

- Turing machines are **deterministic**: for every combination of a (non-accepting, non-rejecting) state  $q$  and a tape symbol  $a \in \Gamma$ , there must be exactly one transition defined for that combination of  $q$  and  $a$ .
- Any transitions that are missing implicitly go straight to a rejecting state. We'll use this later to simplify our designs.



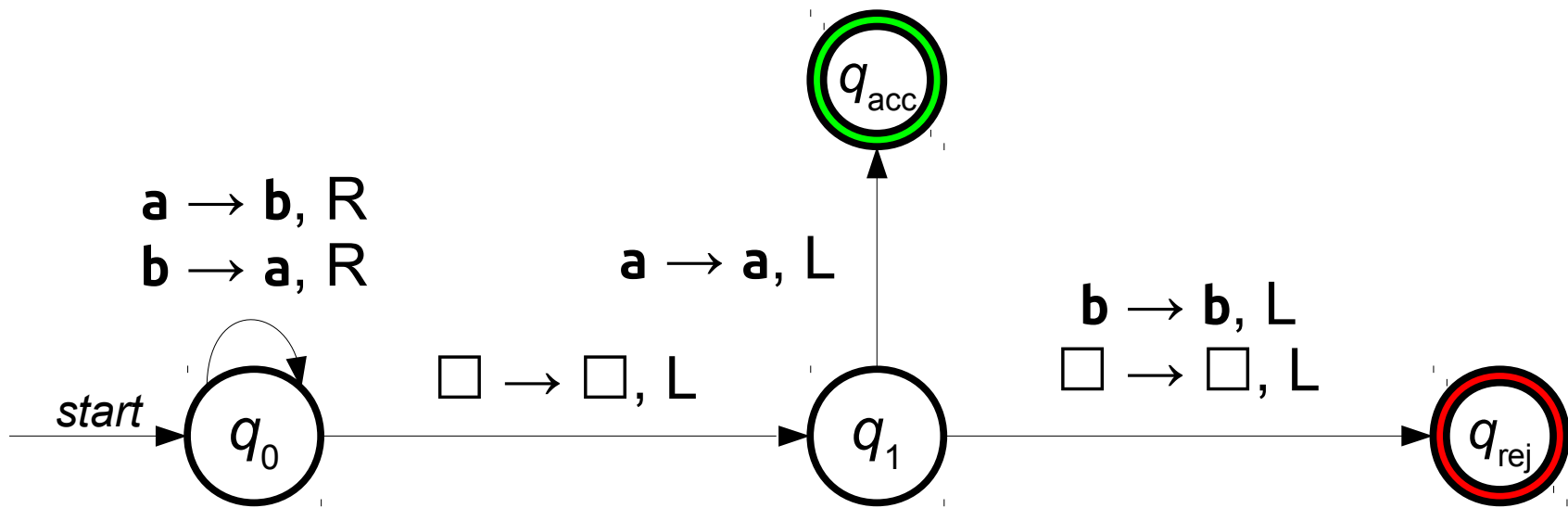
# Determinism

- Turing machines are **deterministic**: for every combination of a (non-accepting, non-rejecting) state  $q$  and a tape symbol  $a \in \Gamma$ , there must be exactly one transition defined for that combination of  $q$  and  $a$ .
- Any transitions that are missing implicitly go straight to a rejecting state. We'll use this later to simplify our designs.



This machine is exactly the same as the previous one.





Run the TM shown above on the input string **bba**.  
 What will the tape look like when the TM finishes running?

**A.** ... b b a ...

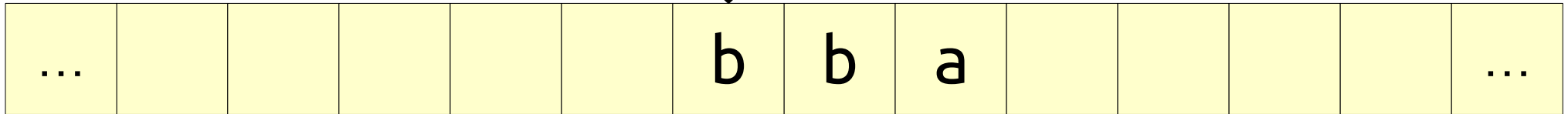
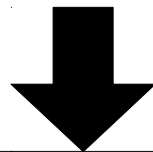
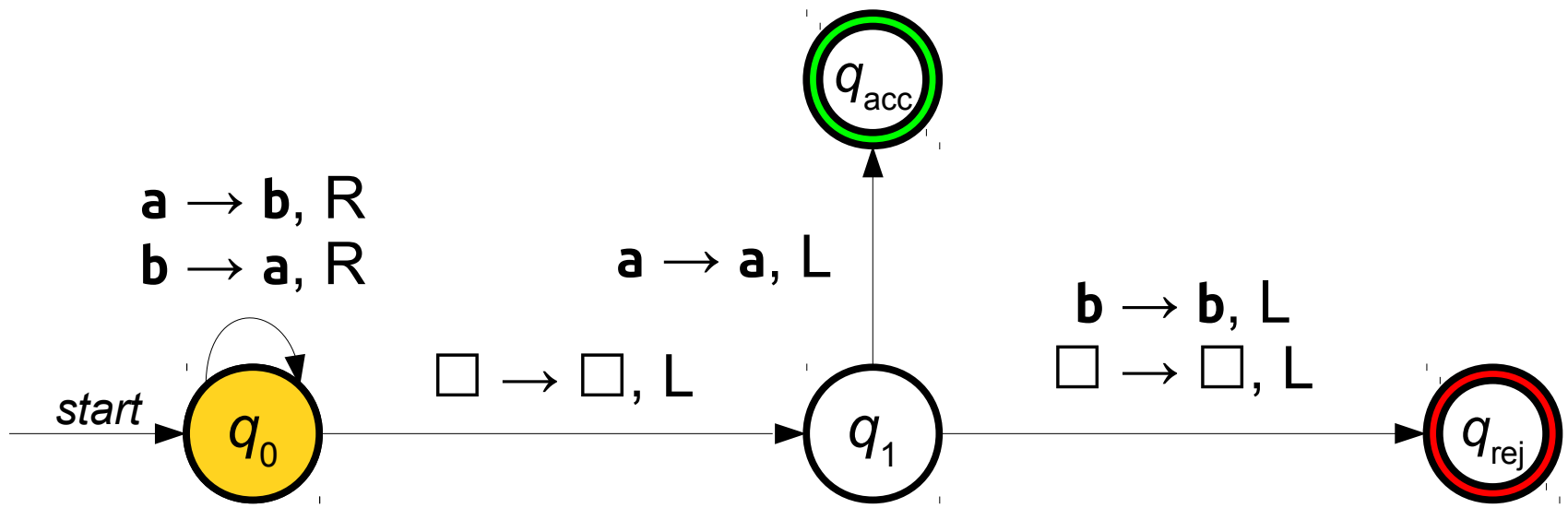
**B.** ... a a b ...

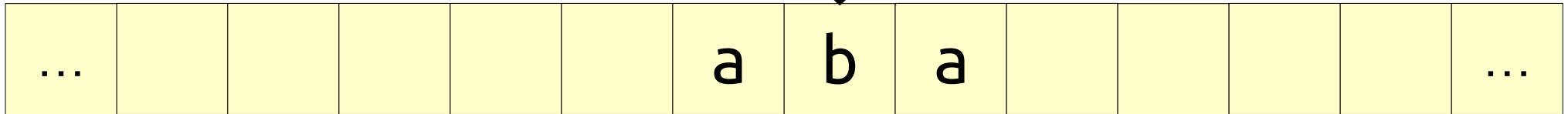
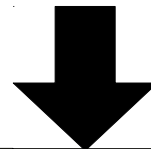
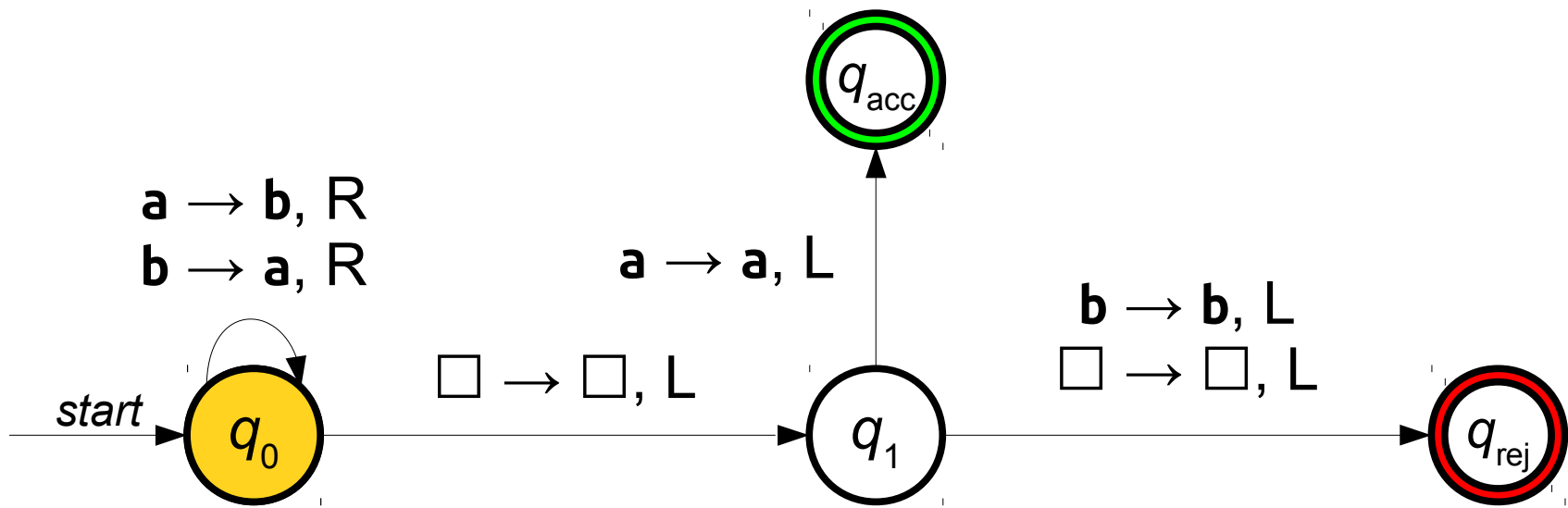
**C.** ... b b a ...

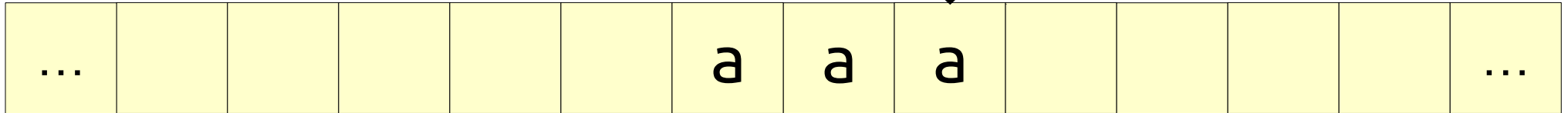
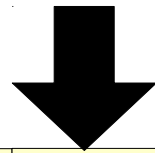
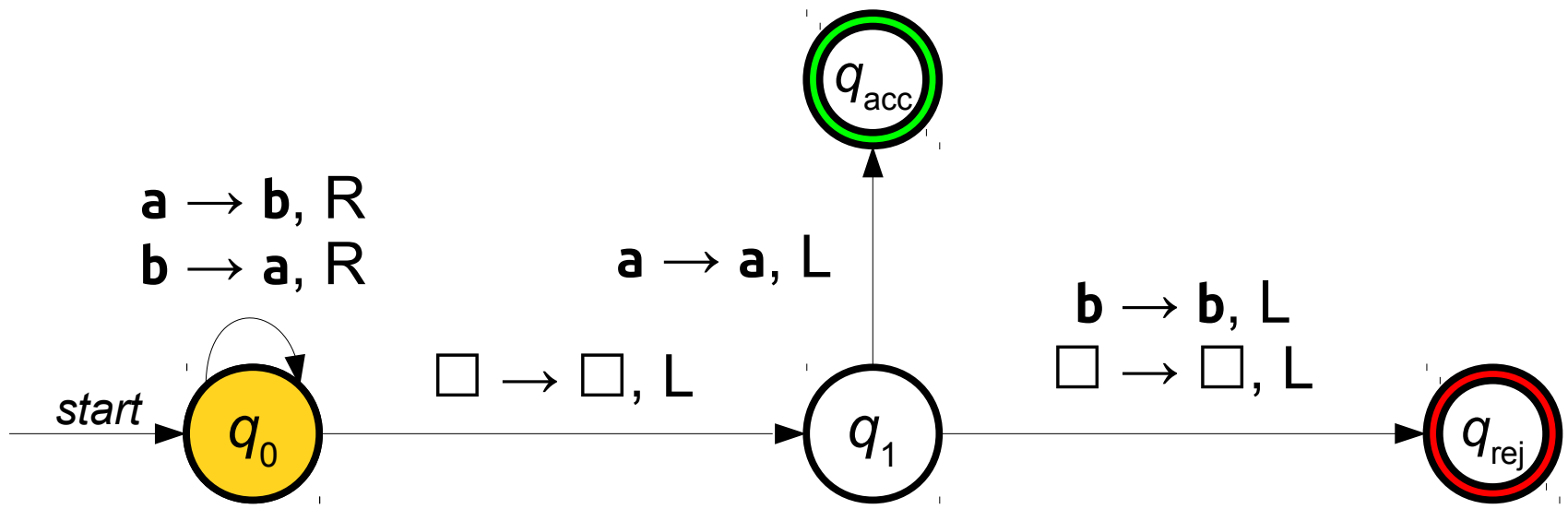
**D.** ... a a b ...

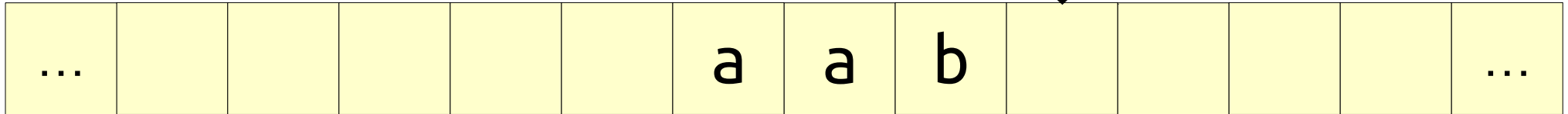
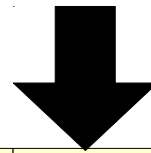
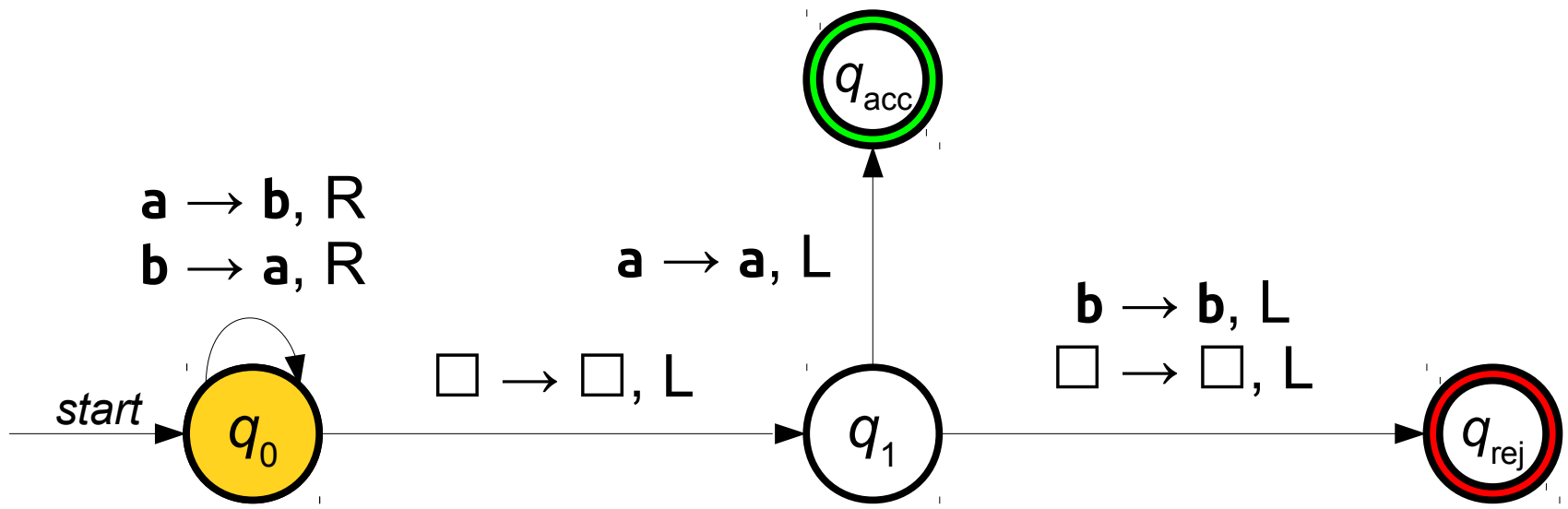
**E.** None of these, or two or more of these.

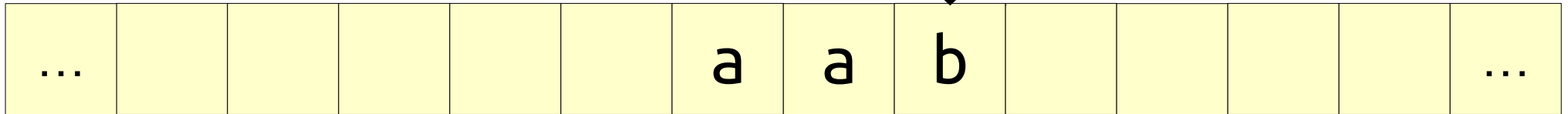
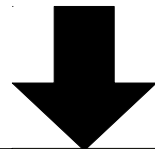
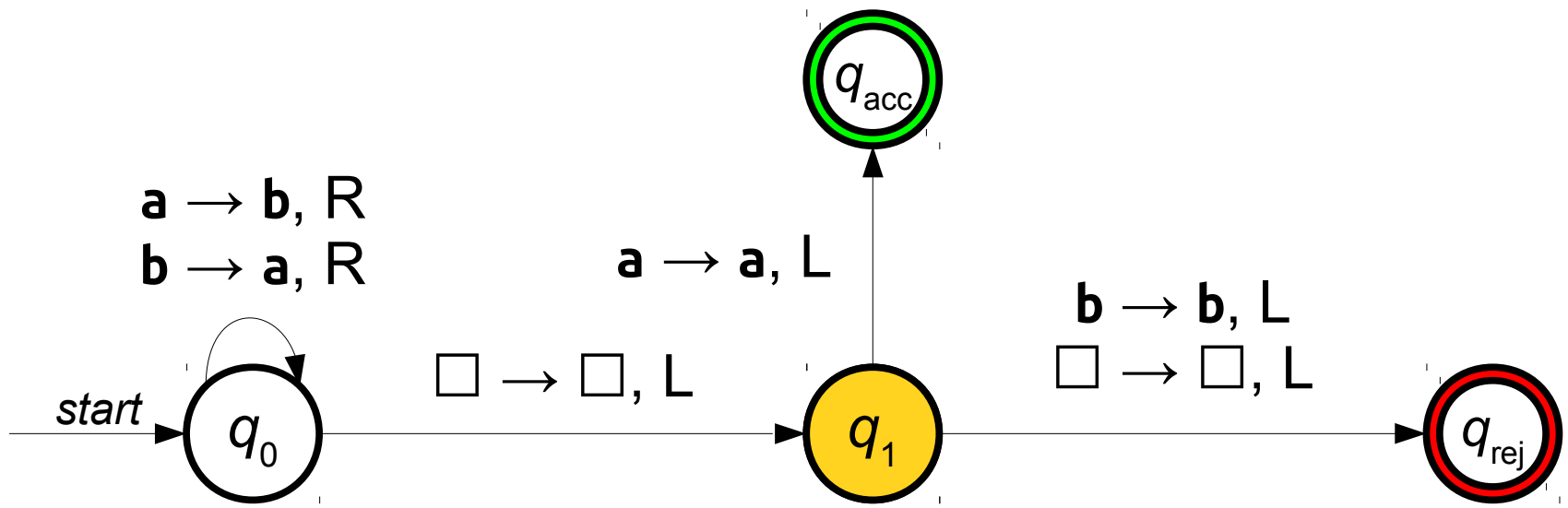
Answer at [Pollevo.com/cs103](https://www.pollevo.com/cs103) or  
 text **CS103** to **22333** once to join, then **A, B, C, D, or E.**

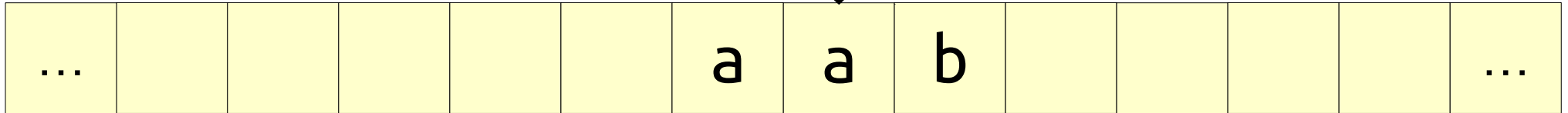
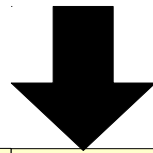
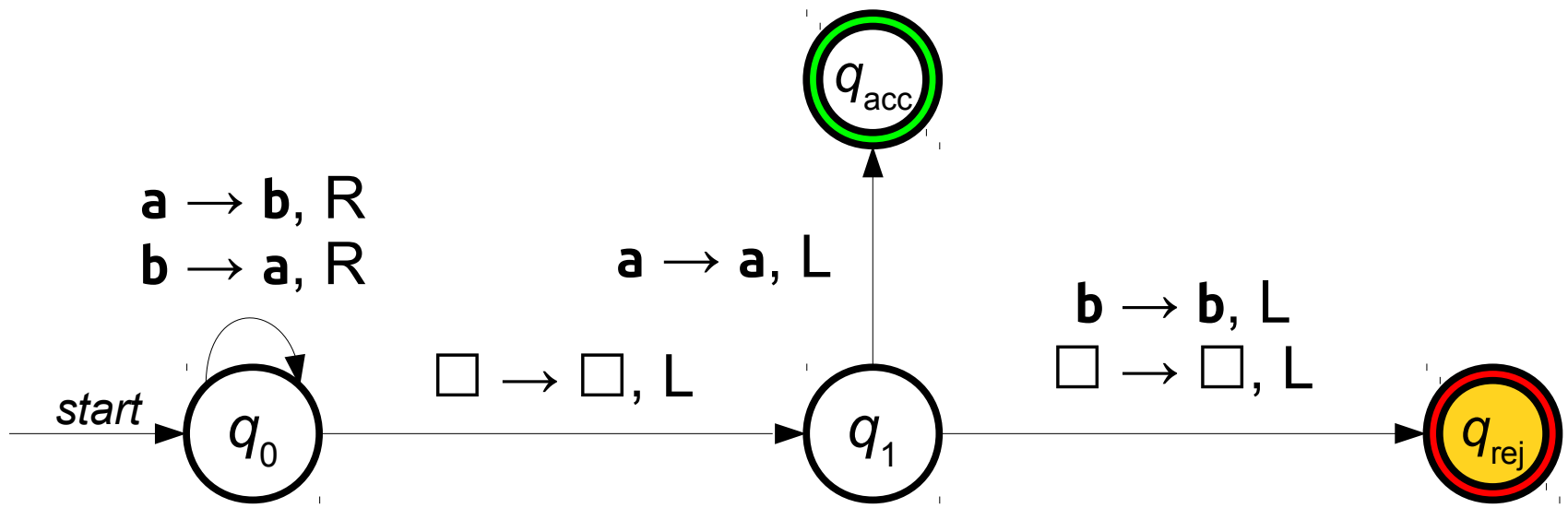


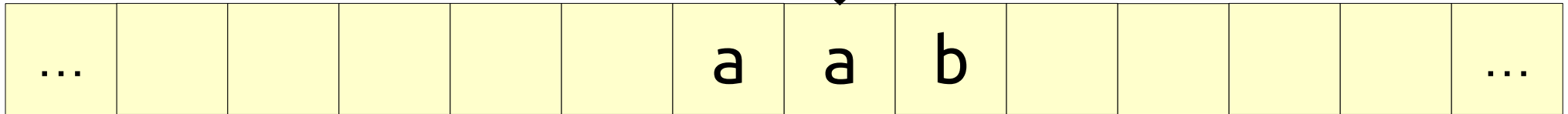
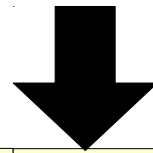
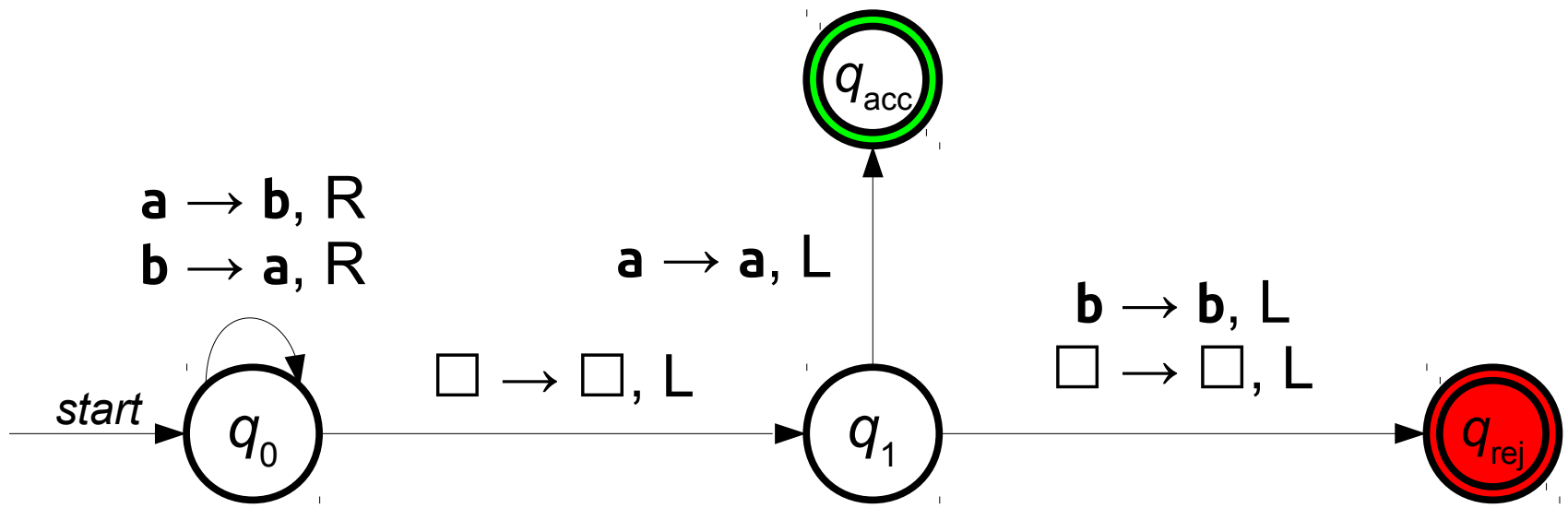




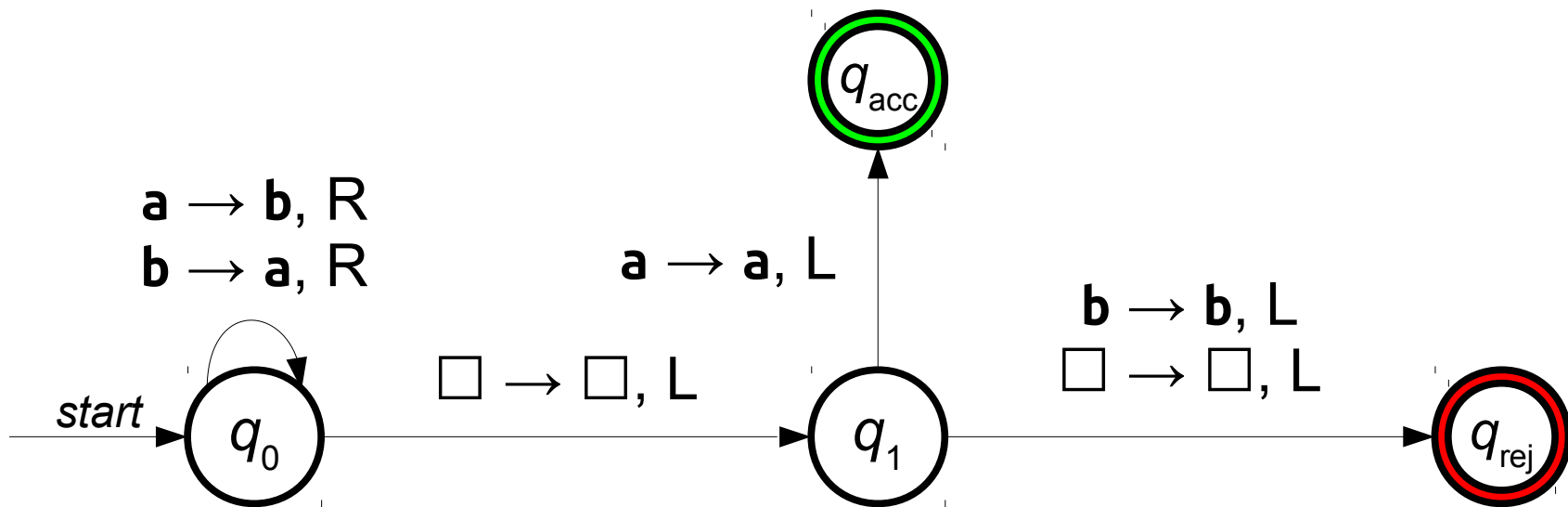












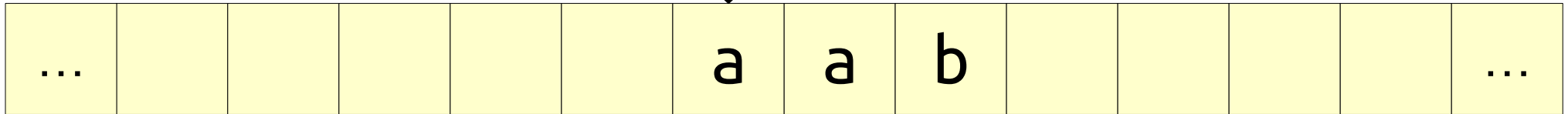
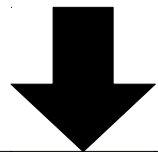
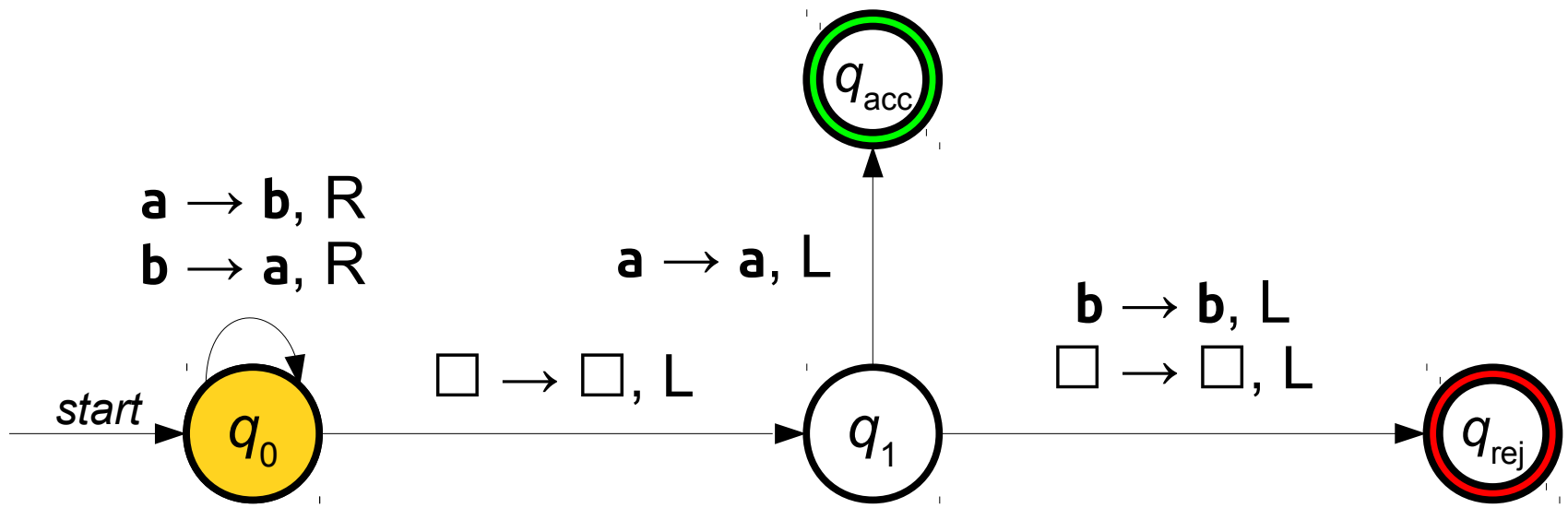
If  $M$  is a Turing machine with input alphabet  $\Sigma$ , then the **language of  $M$** , denoted  $\mathcal{L}(M)$ , is the set

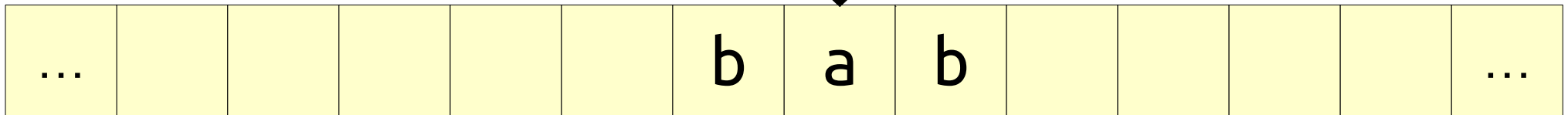
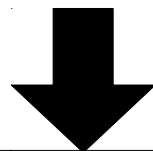
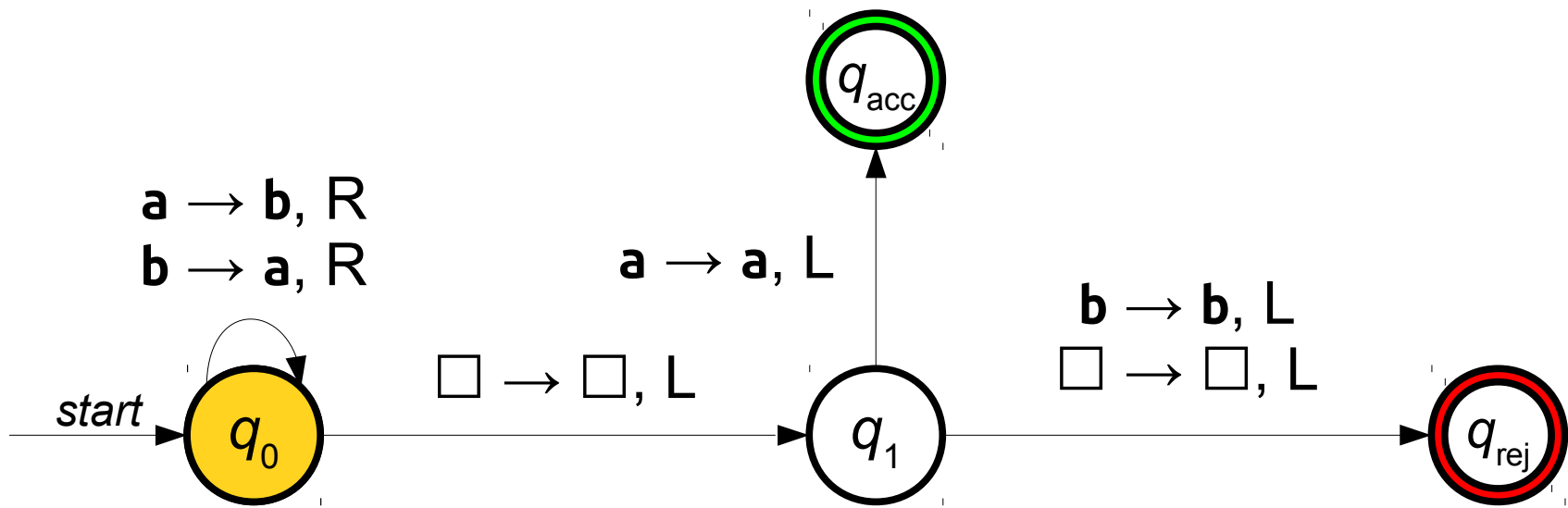
$$\mathcal{L}(M) = \{ w \in \Sigma^* \mid M \text{ accepts } w \}$$

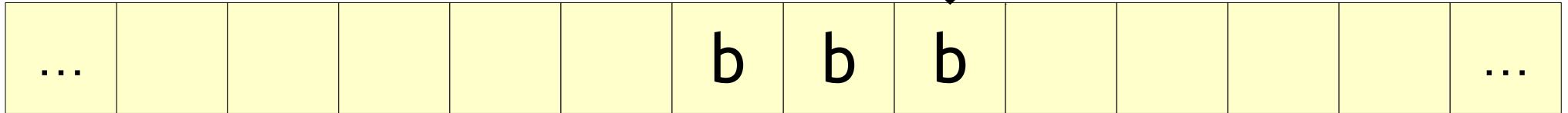
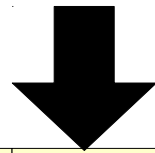
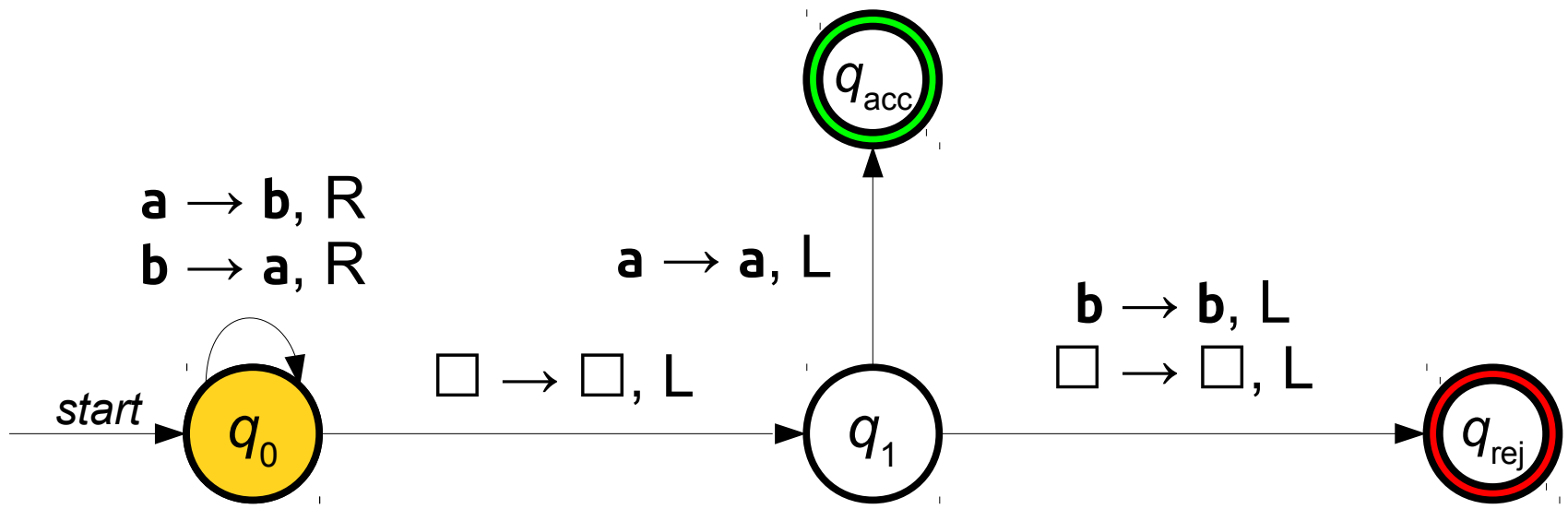
Let  $M$  be the above TM, and assume its input alphabet is  $\{a, b\}$ . What is  $\mathcal{L}(M)$ ?

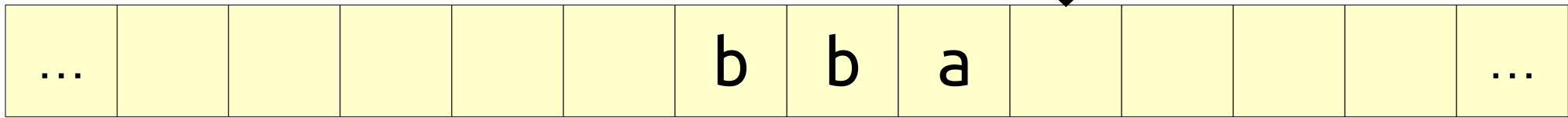
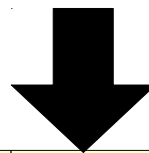
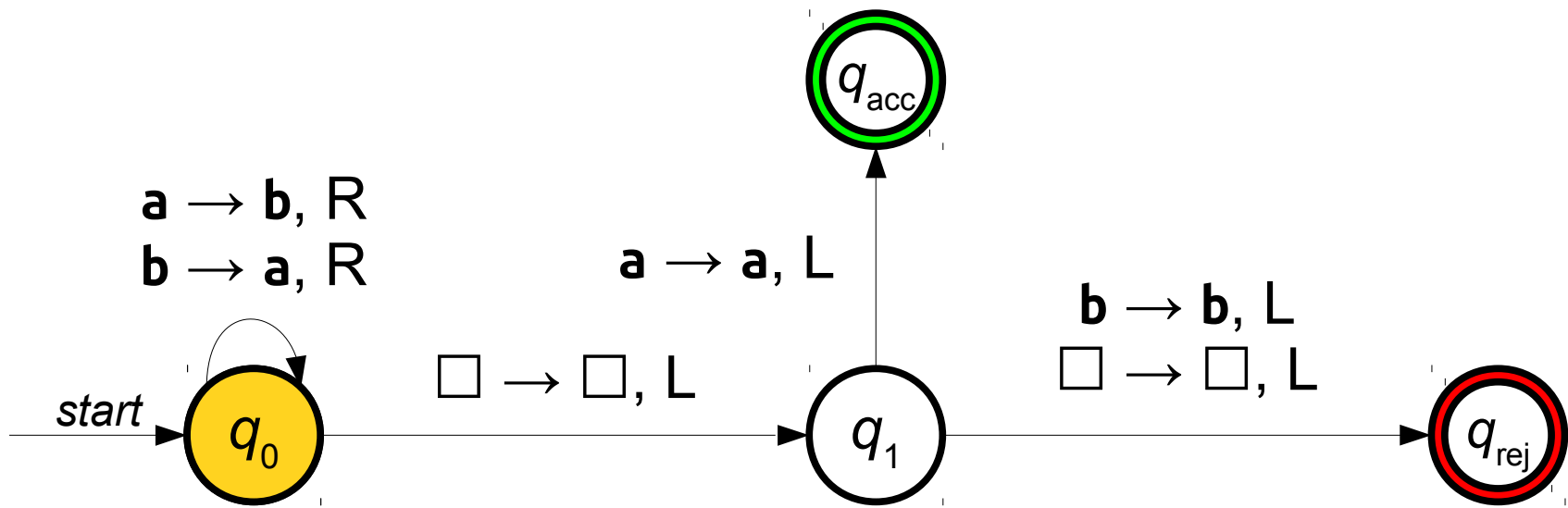
- A.  $\{ w \in \{a, b\}^* \mid w \text{ ends in } a \}$
- B.  $\{ w \in \{a, b\}^* \mid w \text{ ends in } b \}$
- C.  $\emptyset$
- D. None of these, or two or more of these.

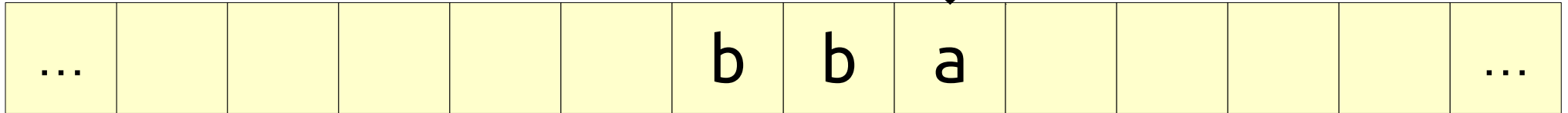
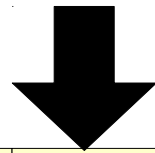
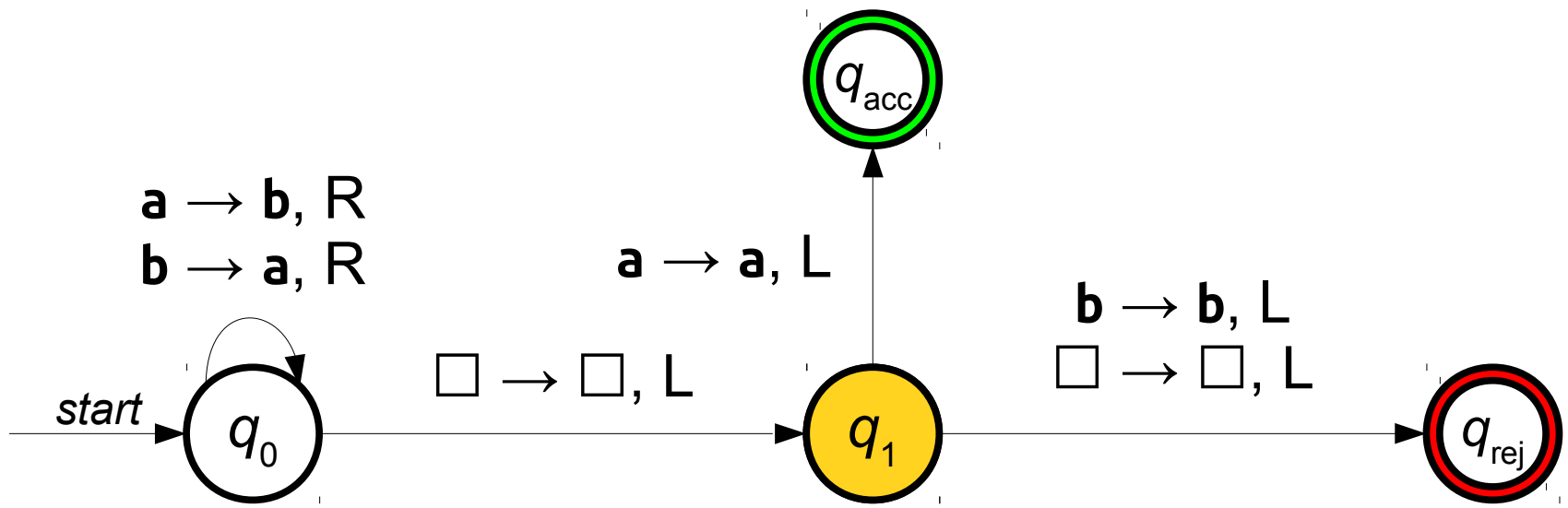
Answer at [PollEv.com/cs103](https://www.pollEv.com/cs103) or  
text **CS103** to **22333** once to join, then **A**, **B**, **C**, or **D**.

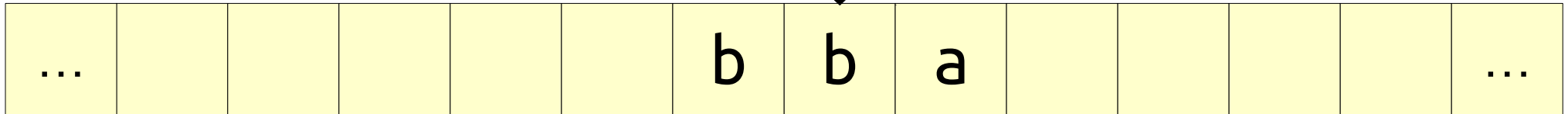
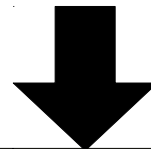
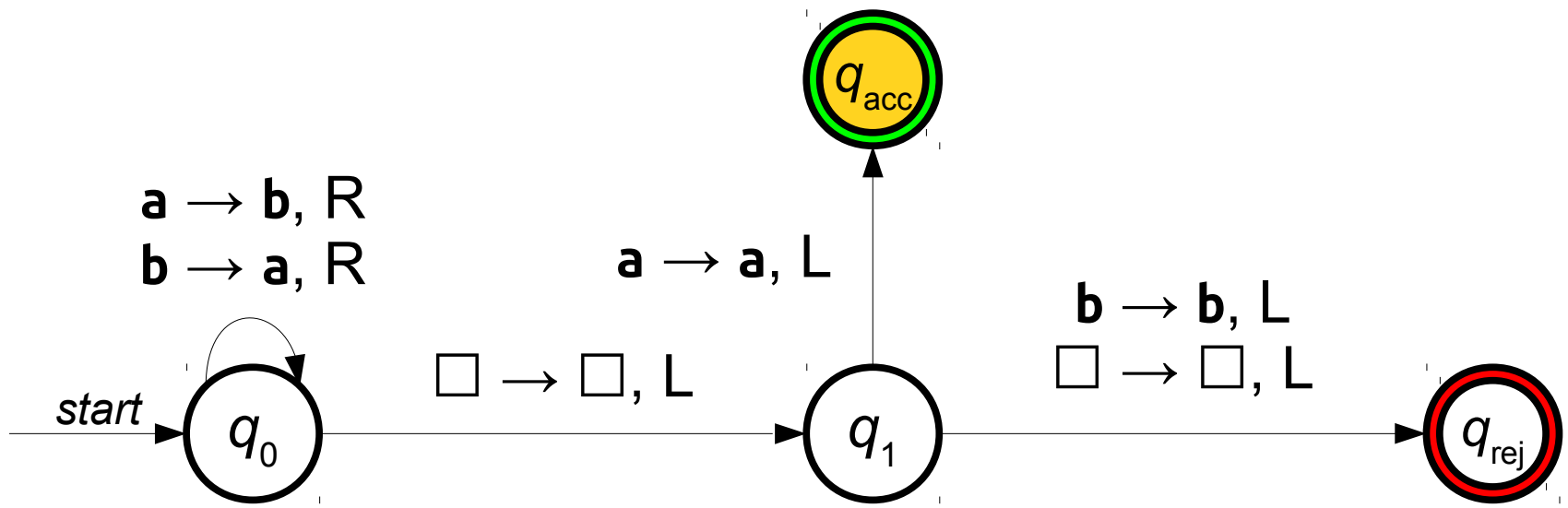


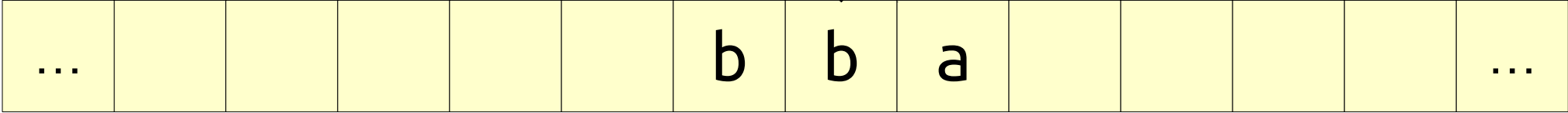
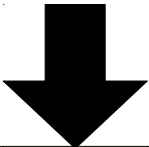
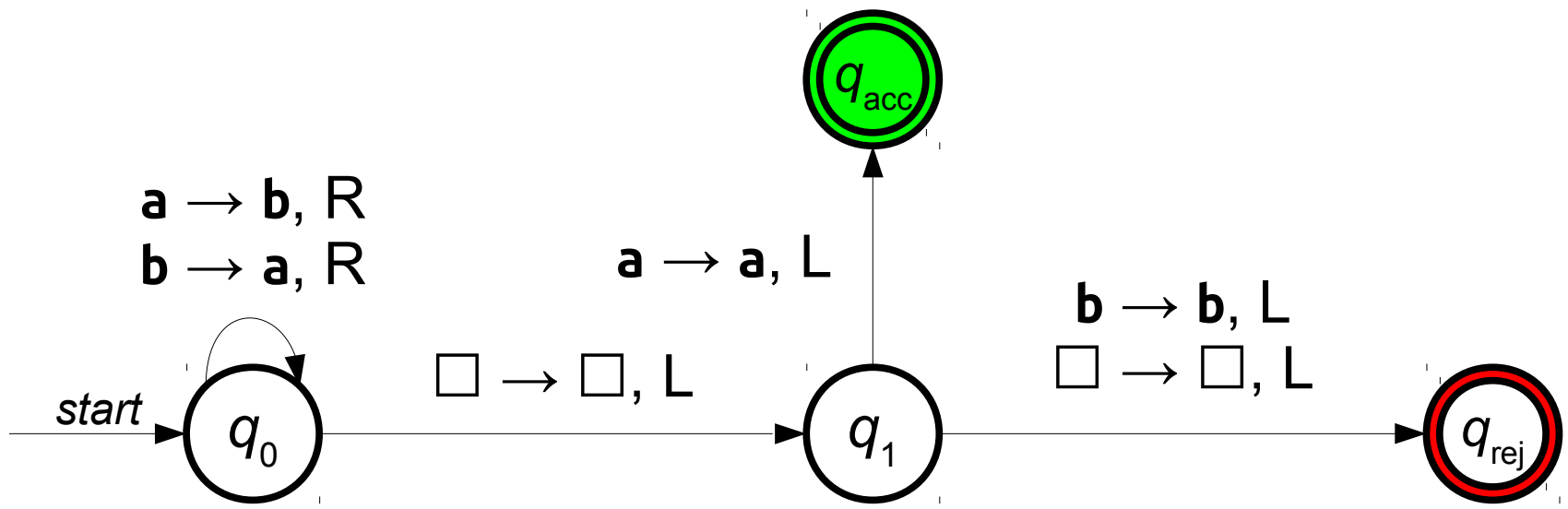




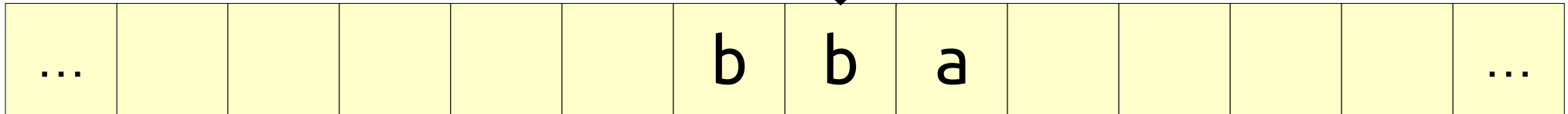
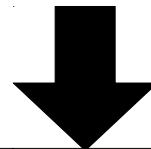
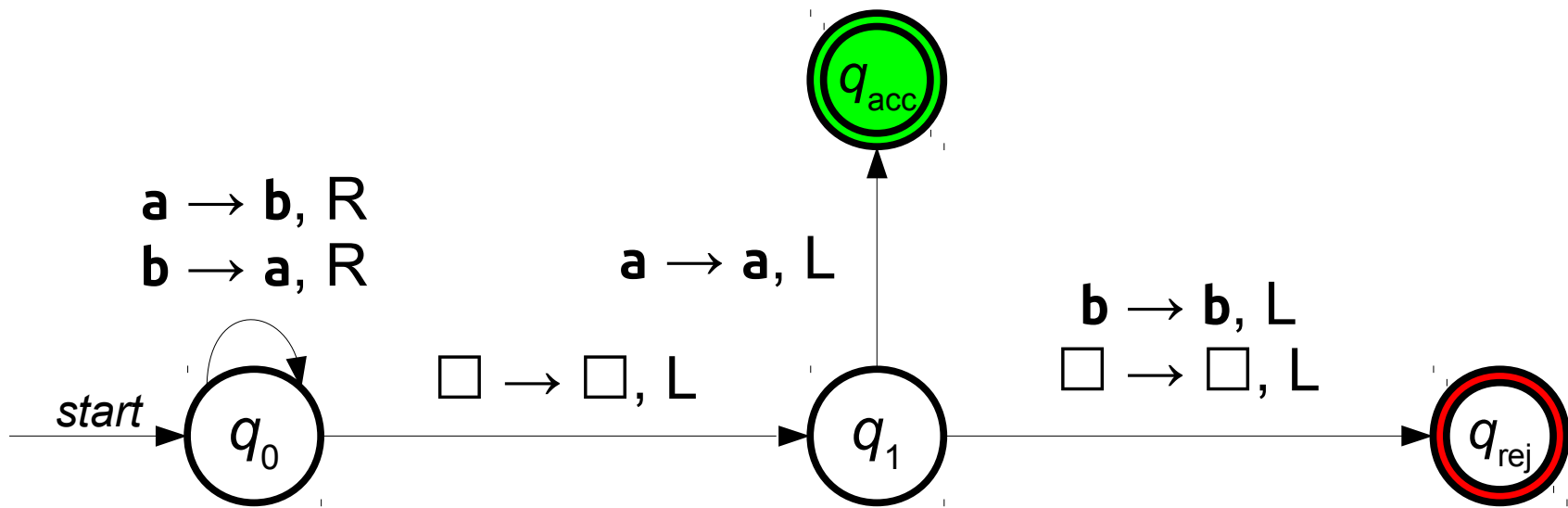












Although the tape ends with **bba** written on it, the original input string was **aab**. This shows that the TM accepts **aab**, not **bba**.

$$\text{So } \mathcal{L}(M) = \{ w \in \{a, b\}^* \mid w \text{ ends in } b \}$$

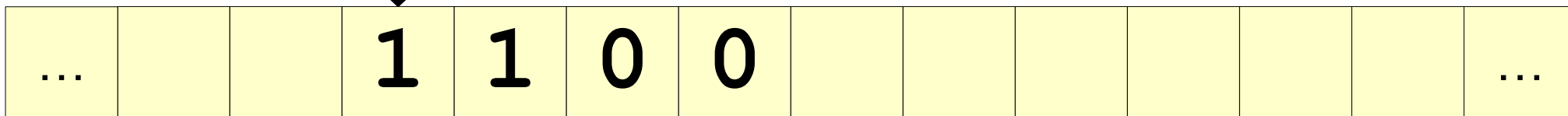
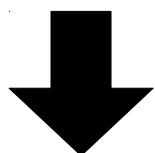
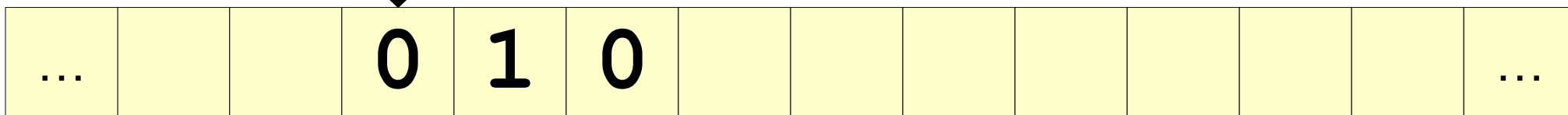
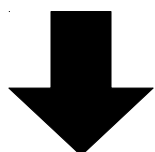
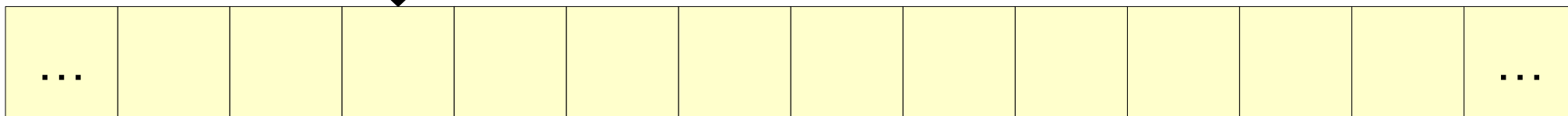
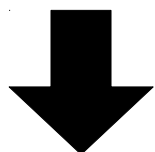
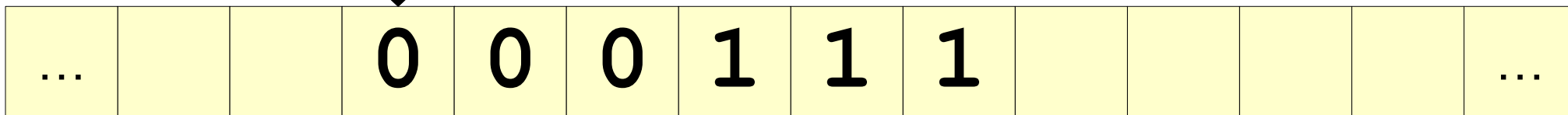
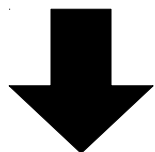
# Designing Turing Machines

- Despite their simplicity, Turing machines are very powerful computing devices.
- Today's lecture explores how to design Turing machines for various languages.

# Designing Turing Machines

- Let  $\Sigma = \{0, 1\}$  and consider the language  $L = \{0^n 1^n \mid n \in \mathbb{N}\}$ .
- We know that  $L$  is context-free.
- How might we build a Turing machine for it?

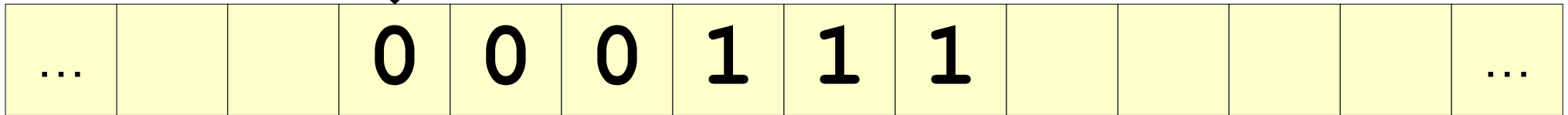
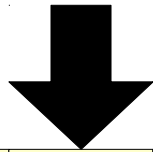
$$L = \{0^n 1^n \mid n \in \mathbb{N}\}$$



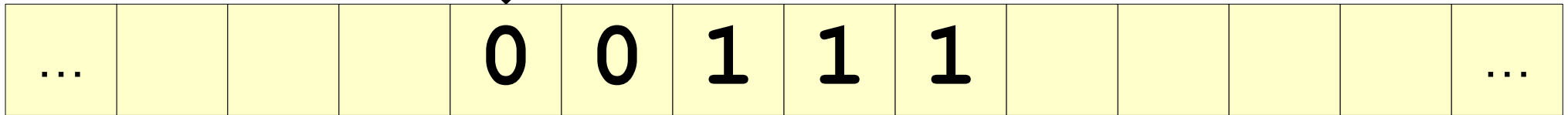
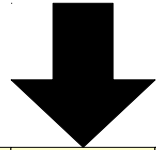
# A Recursive Approach

- The string  $\varepsilon$  is in  $L$ .
- The string  $0w1$  is in  $L$  iff  $w$  is in  $L$ .
- Any string starting with  $1$  is not in  $L$ .
- Any string ending with  $0$  is not in  $L$ .

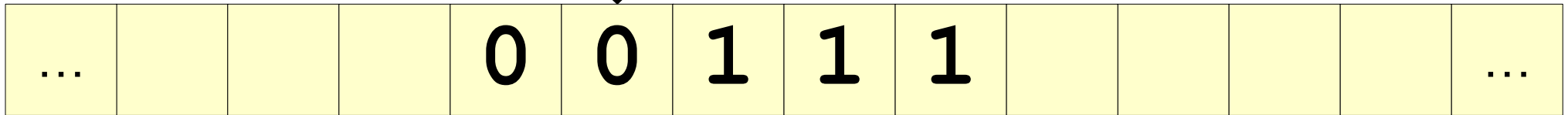
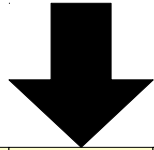
# A Sketch of the TM



# A Sketch of the TM

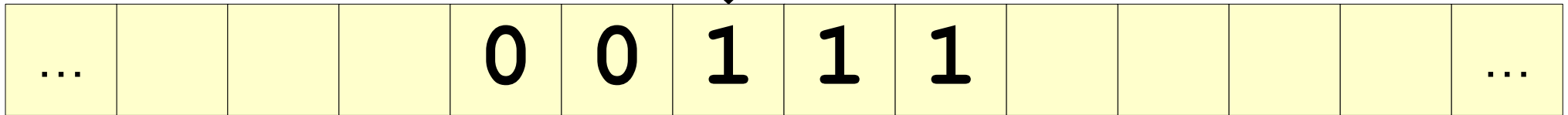
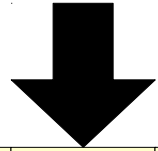


# A Sketch of the TM

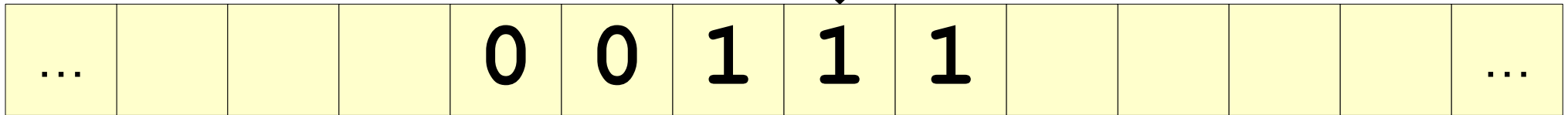
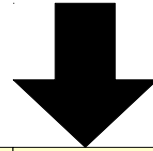




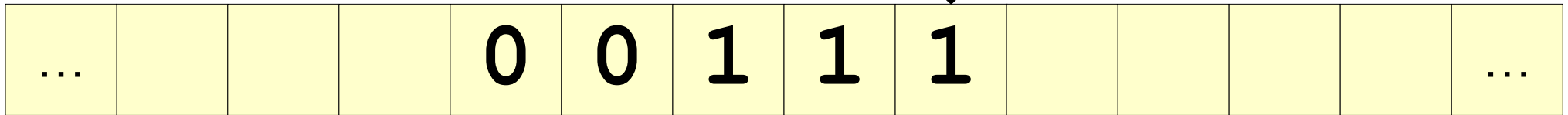
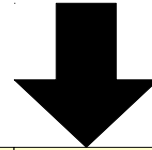
# A Sketch of the TM



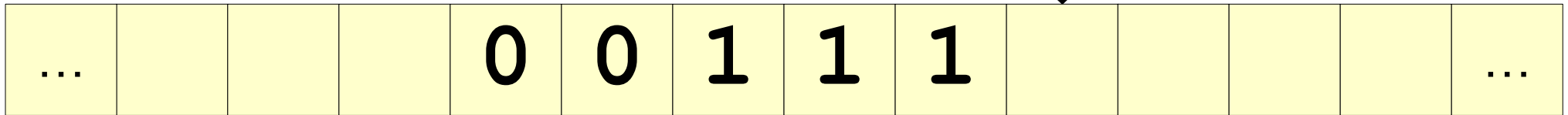
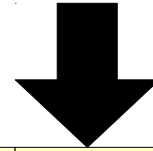
# A Sketch of the TM



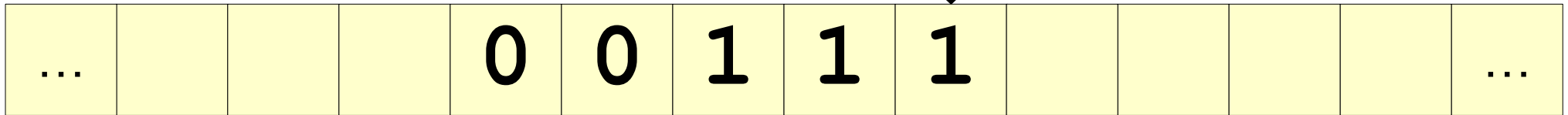
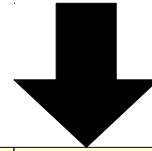
# A Sketch of the TM



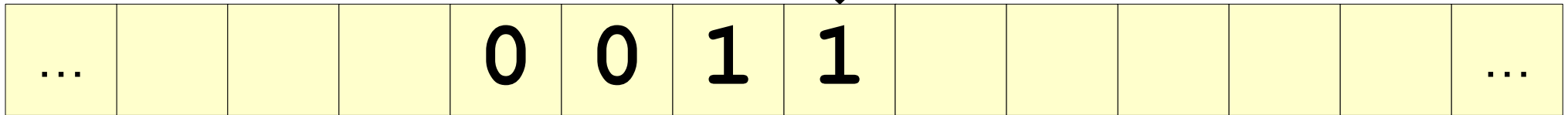
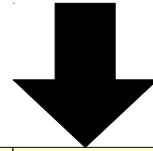
# A Sketch of the TM



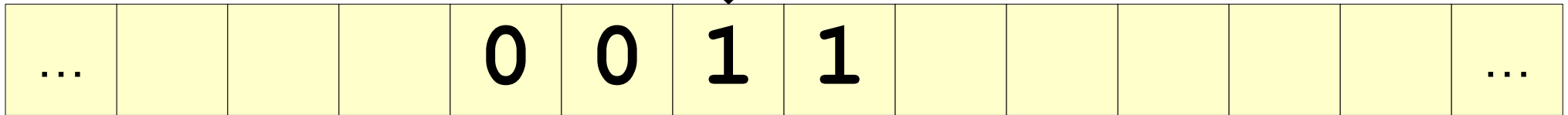
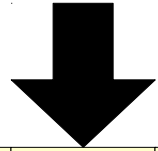
# A Sketch of the TM



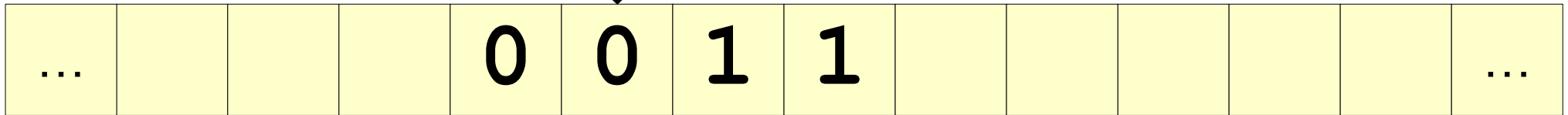
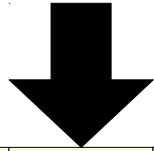
# A Sketch of the TM



# A Sketch of the TM

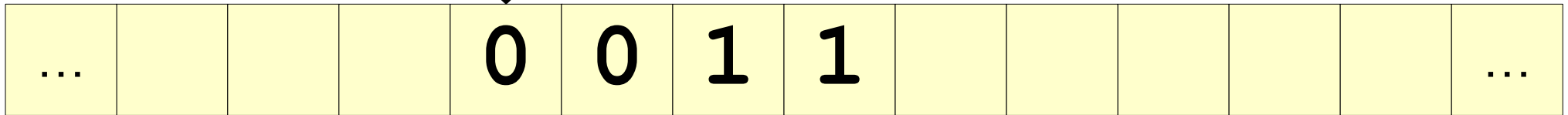
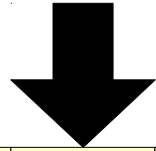


# A Sketch of the TM

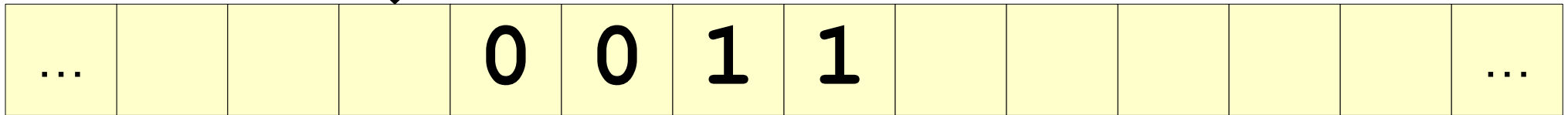
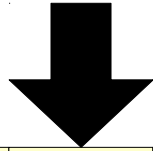




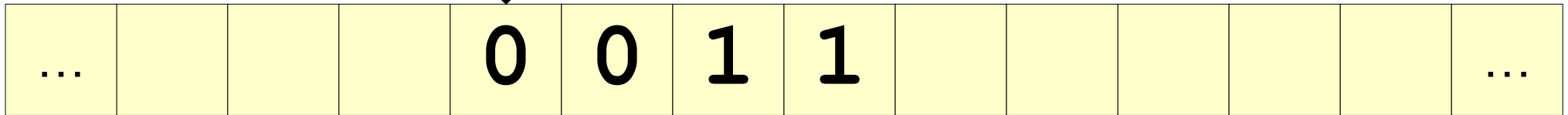
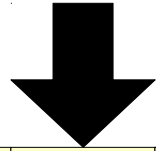
# A Sketch of the TM



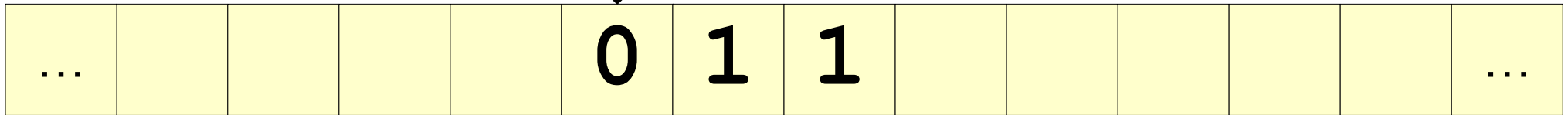
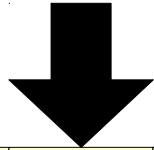
# A Sketch of the TM



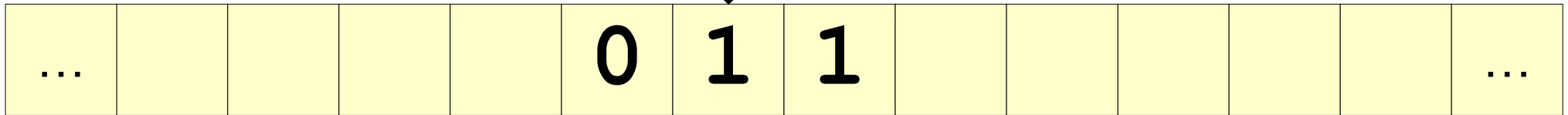
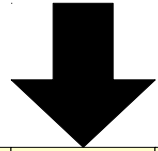
# A Sketch of the TM



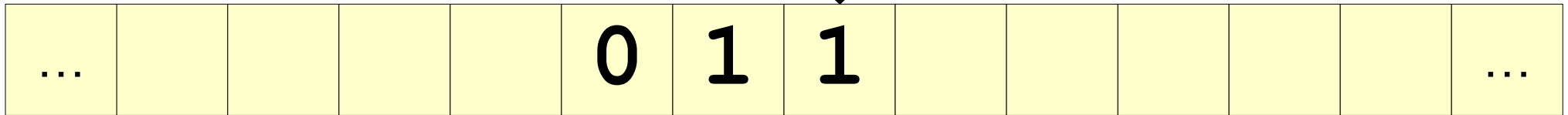
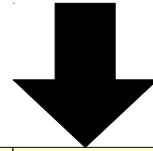
# A Sketch of the TM



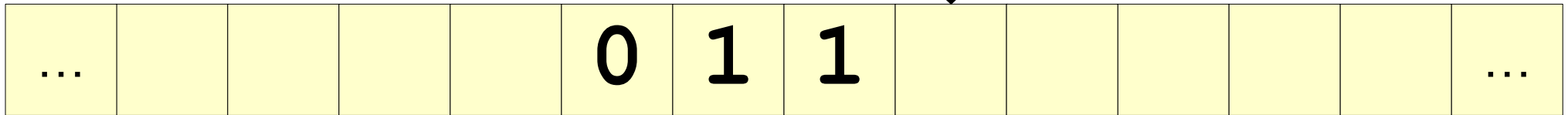
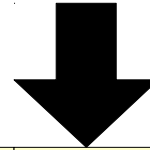
# A Sketch of the TM



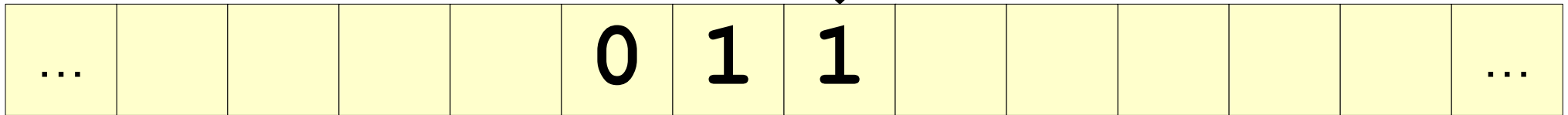
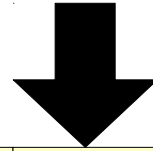
# A Sketch of the TM



# A Sketch of the TM

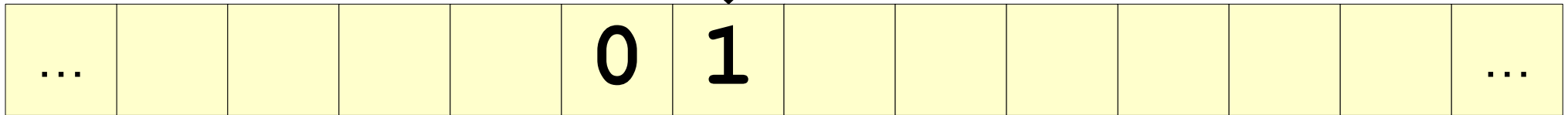
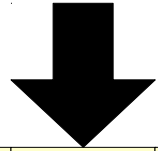


# A Sketch of the TM

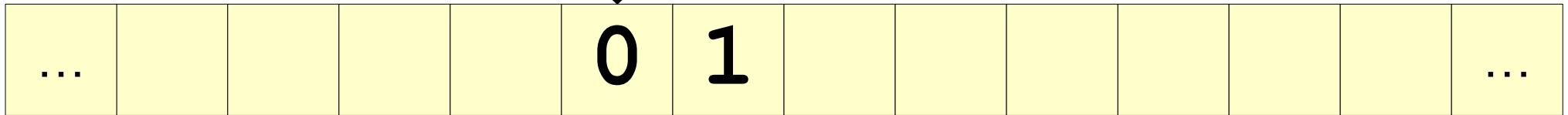
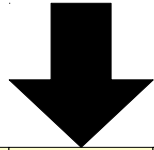




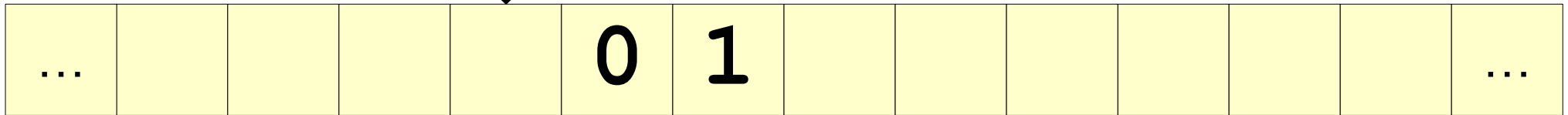
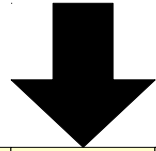
# A Sketch of the TM



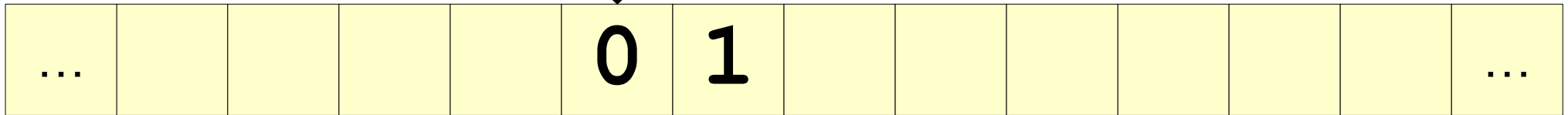
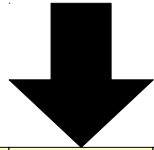
# A Sketch of the TM



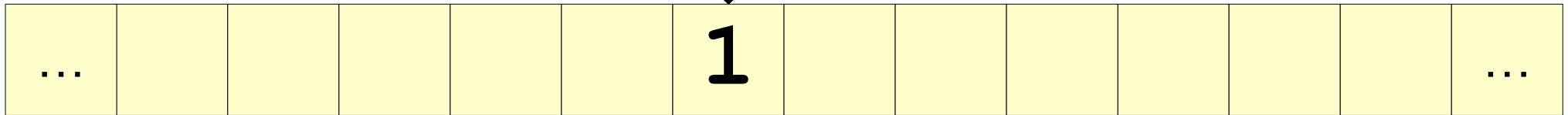
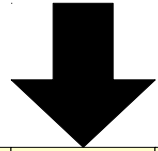
# A Sketch of the TM



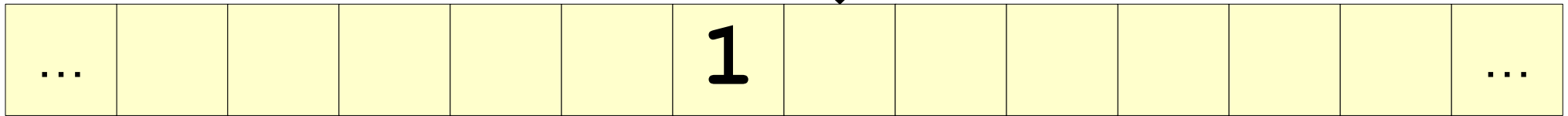
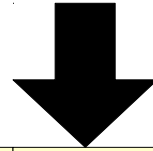
# A Sketch of the TM



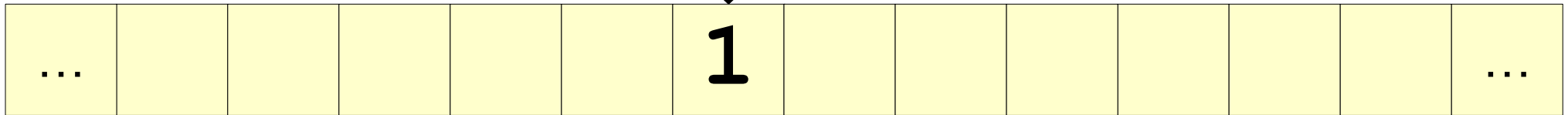
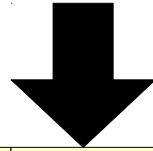
# A Sketch of the TM



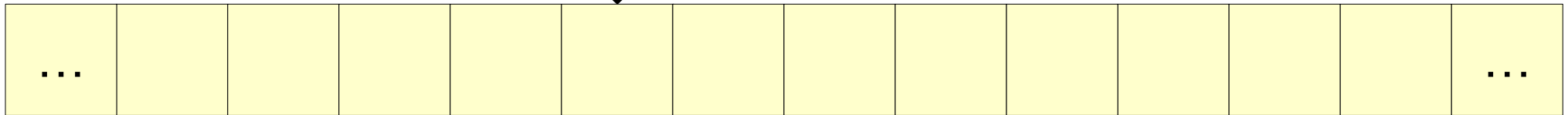
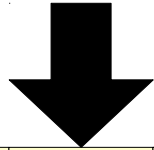
# A Sketch of the TM



# A Sketch of the TM

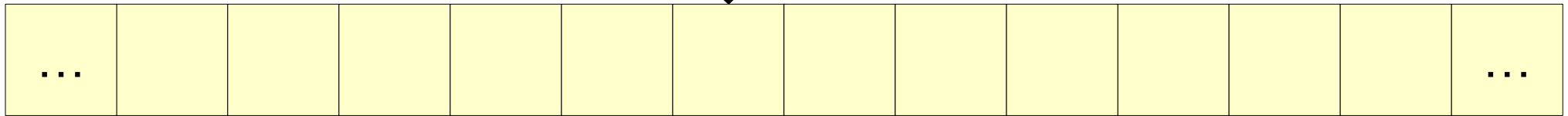
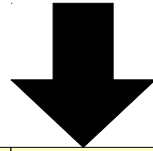


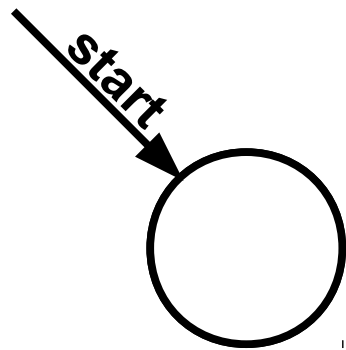
# A Sketch of the TM

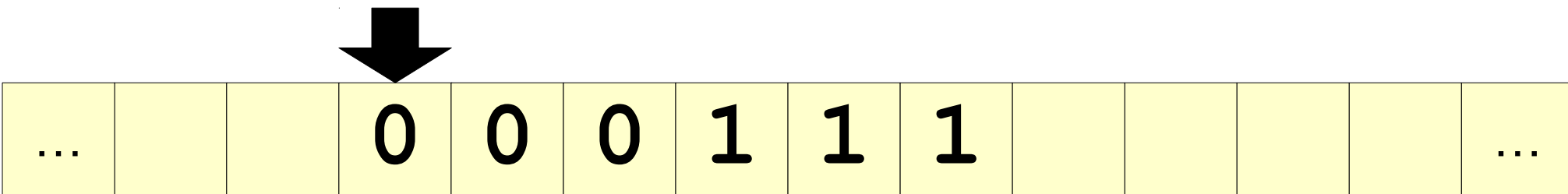
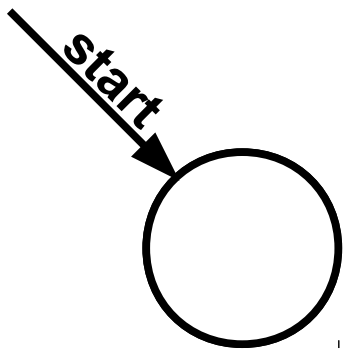


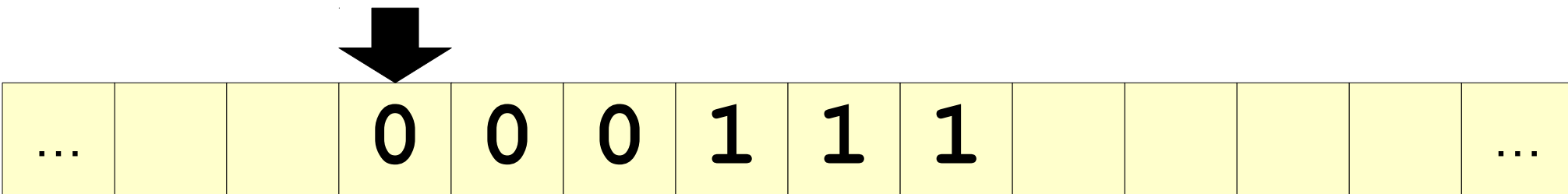
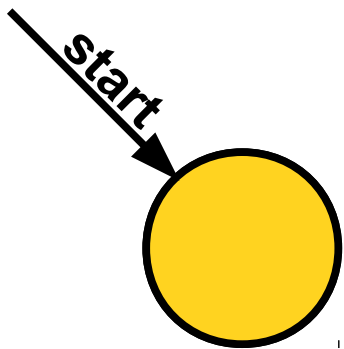


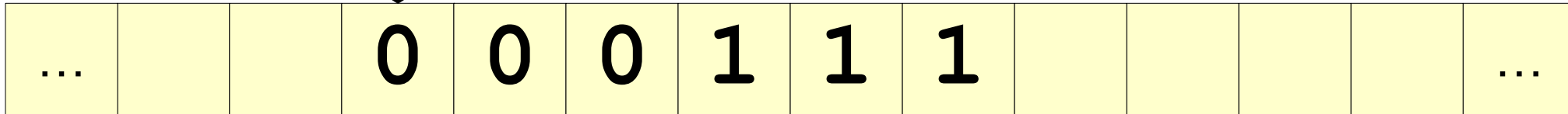
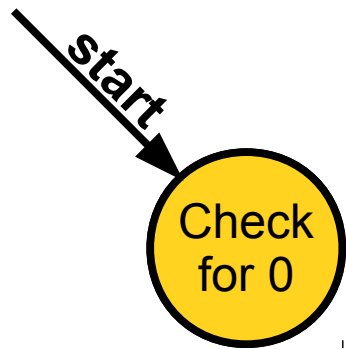
# A Sketch of the TM

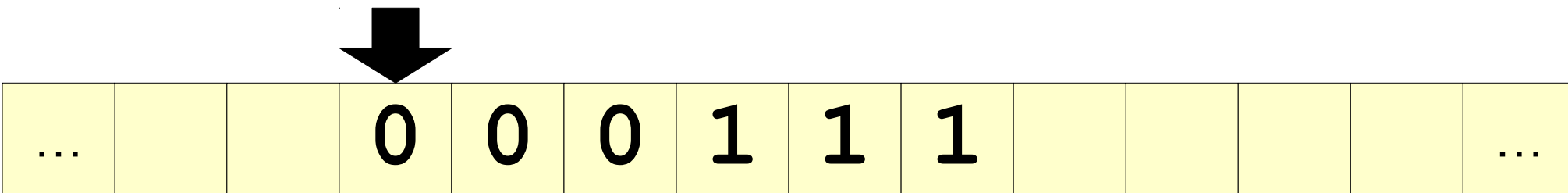
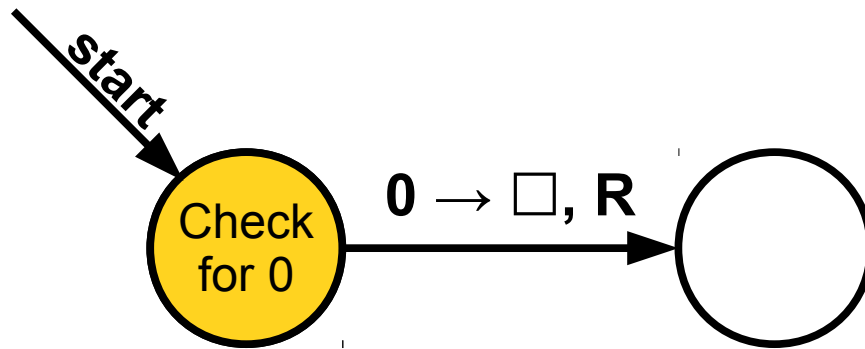


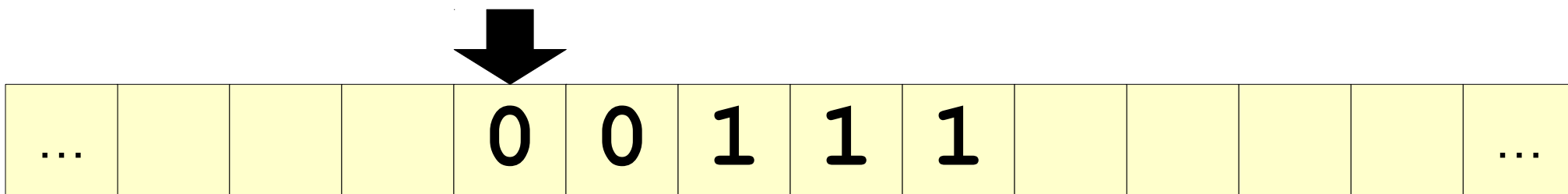
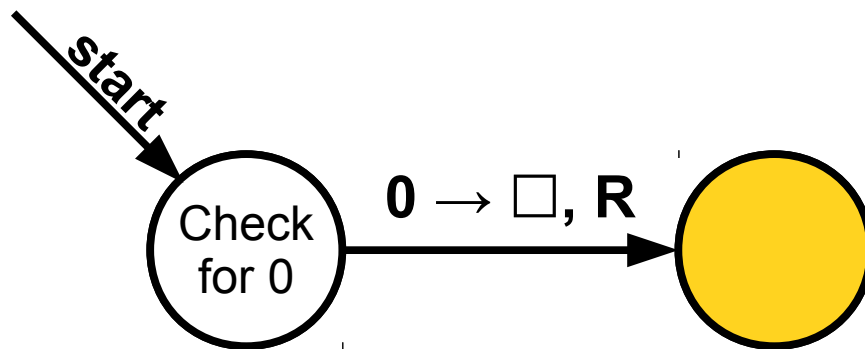


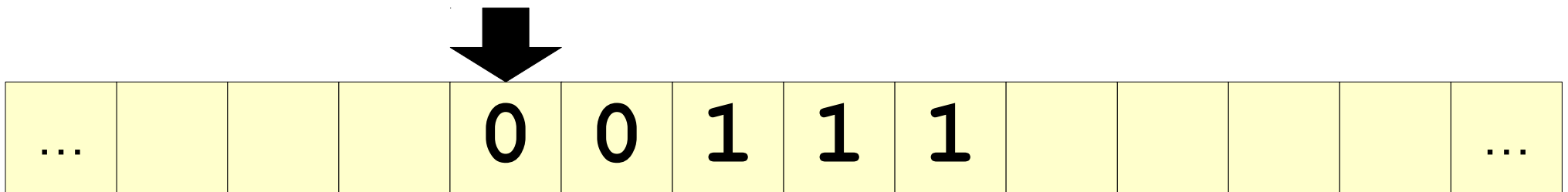
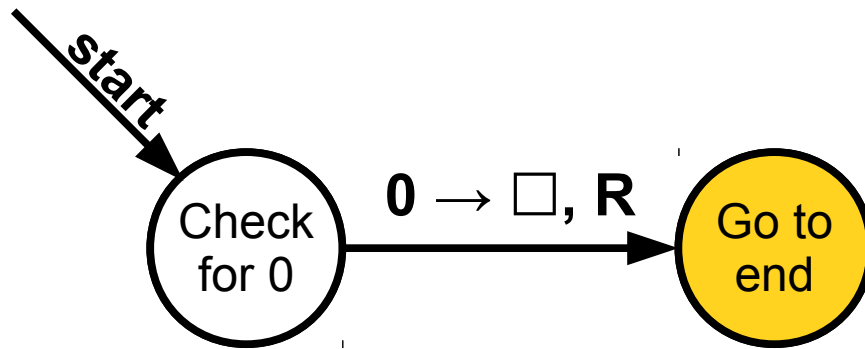




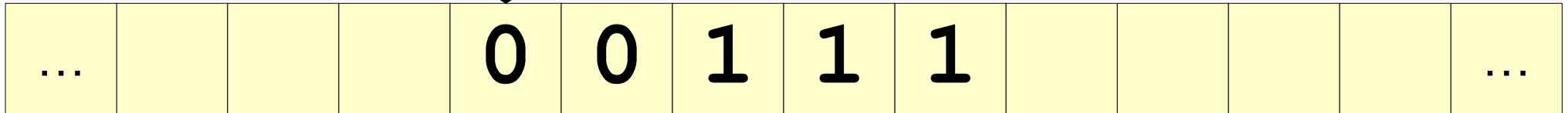
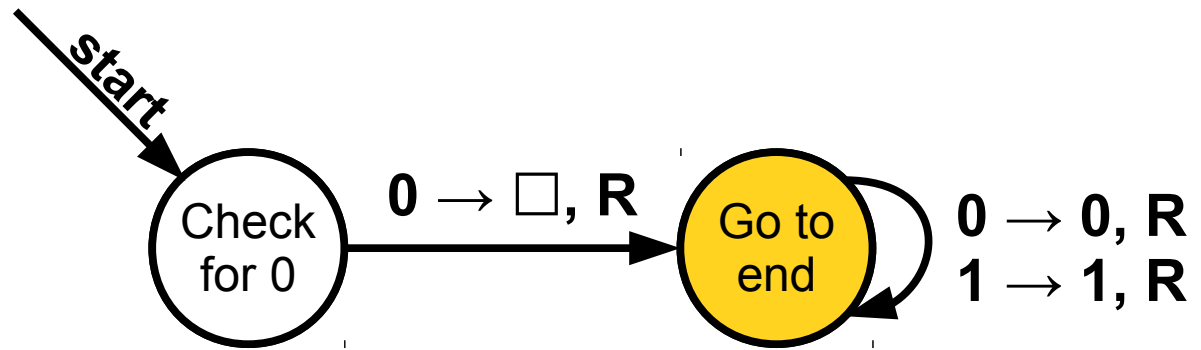


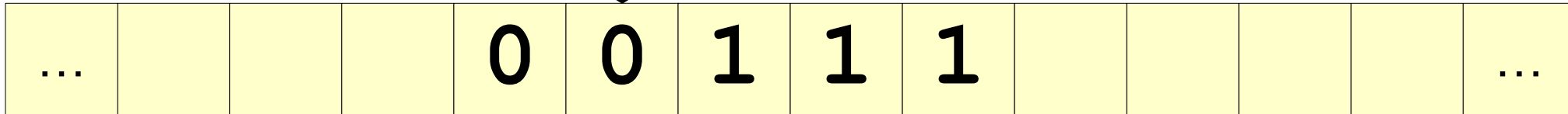
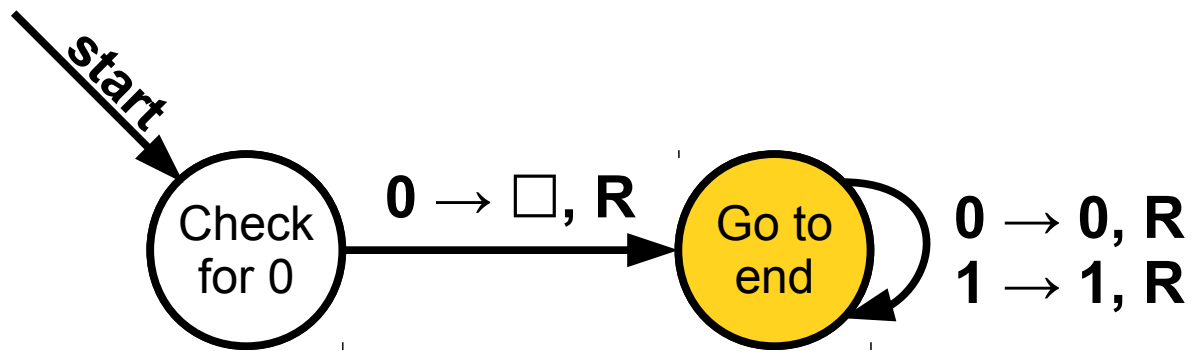


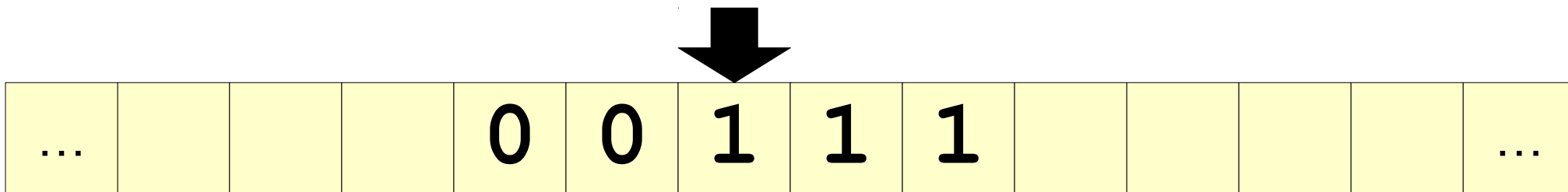
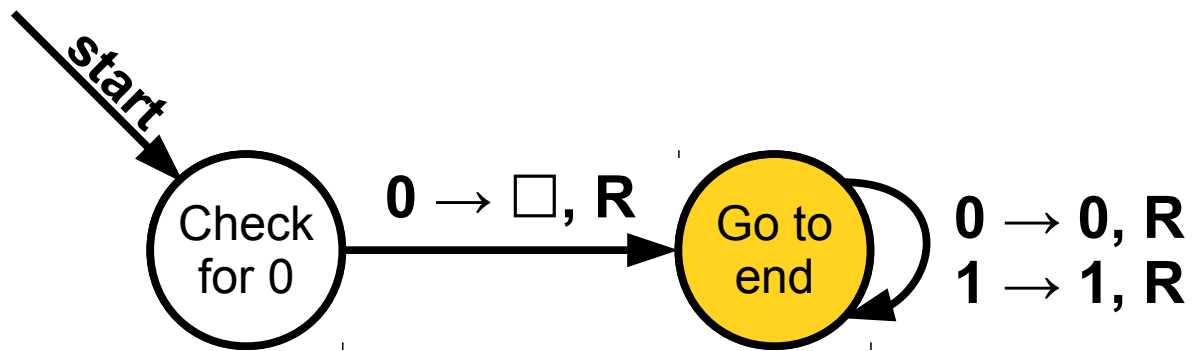


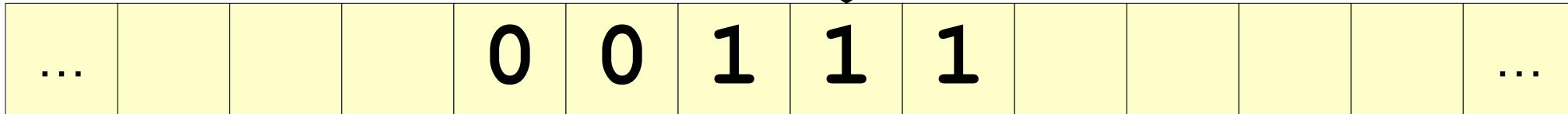
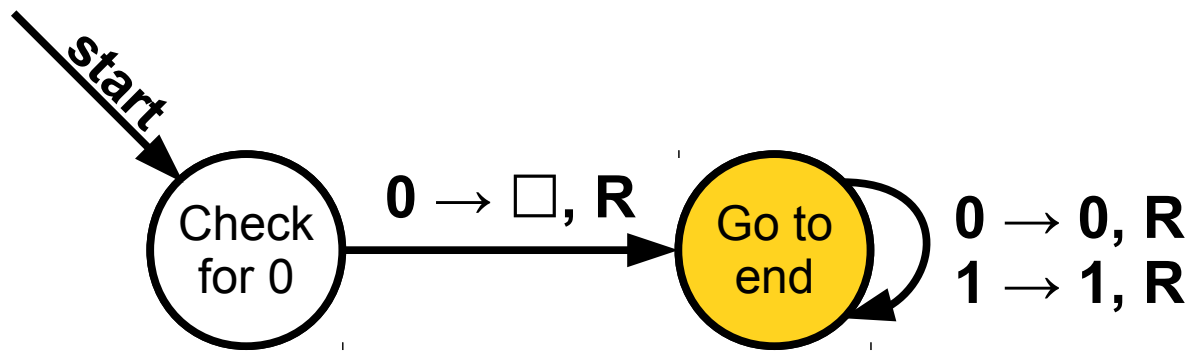


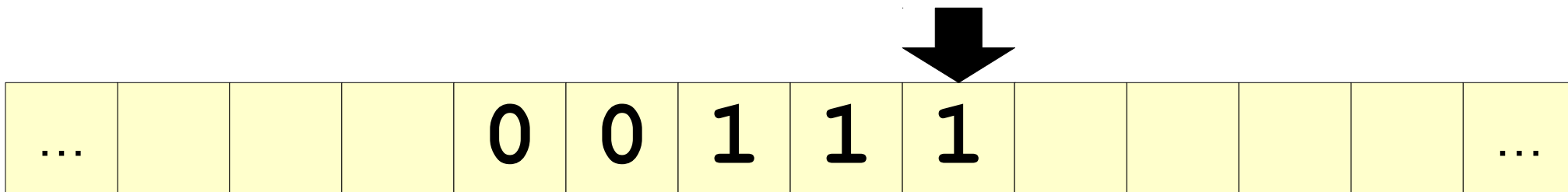
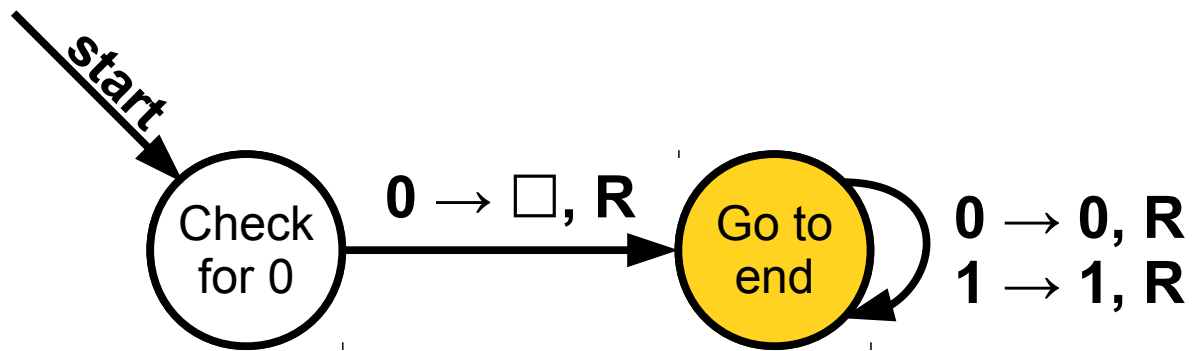


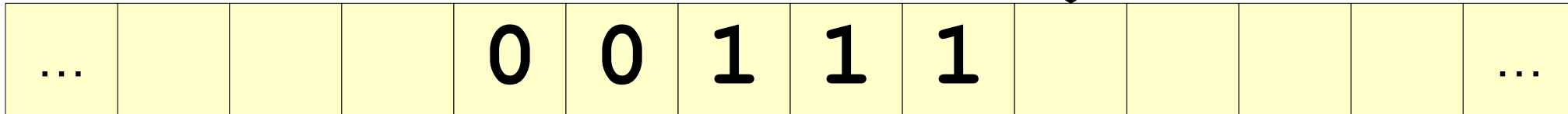
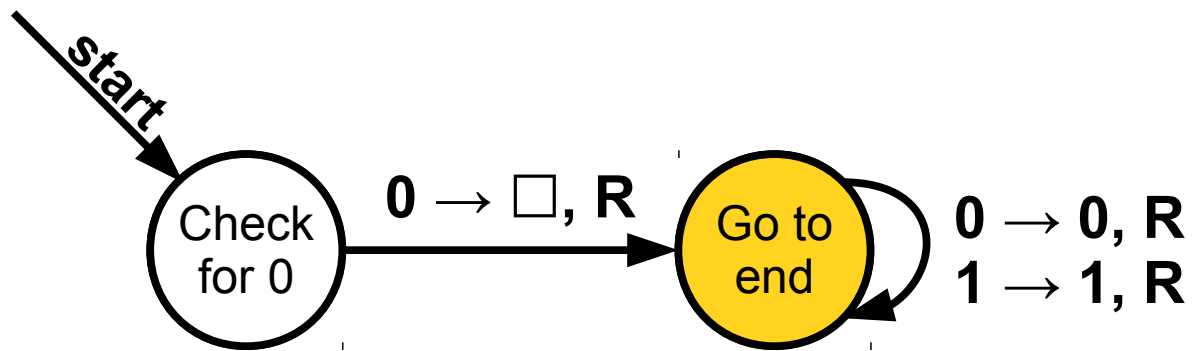


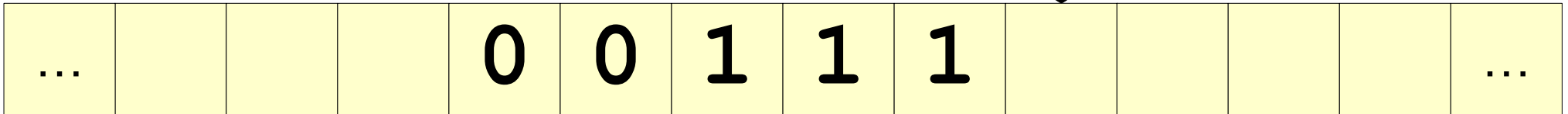
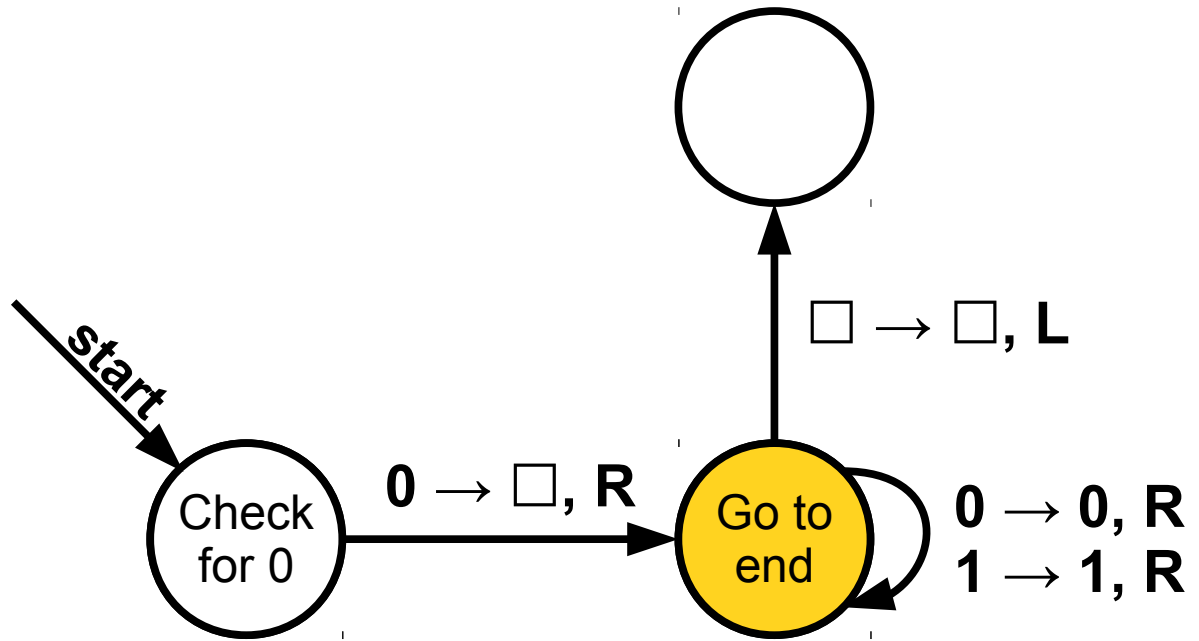


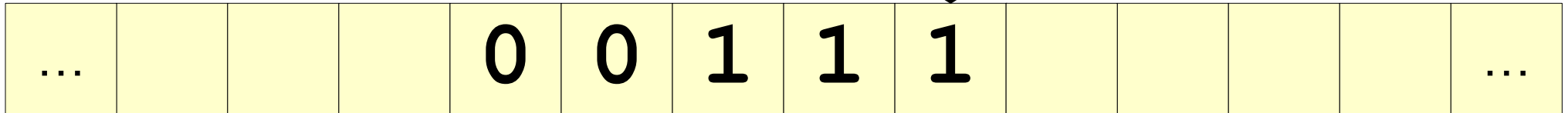
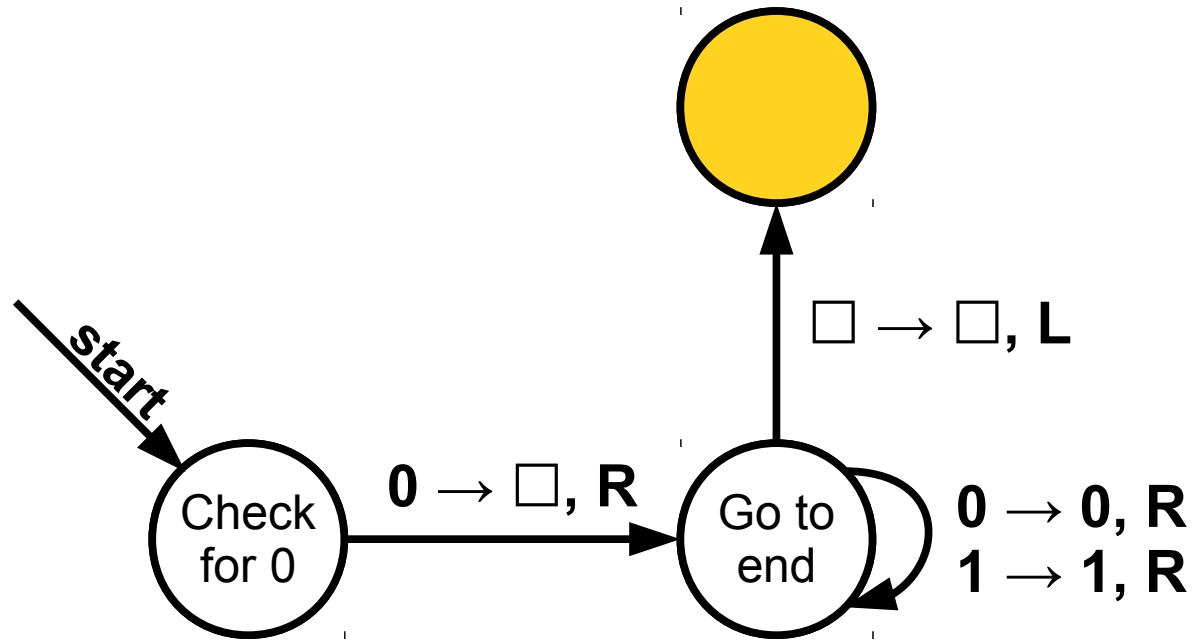




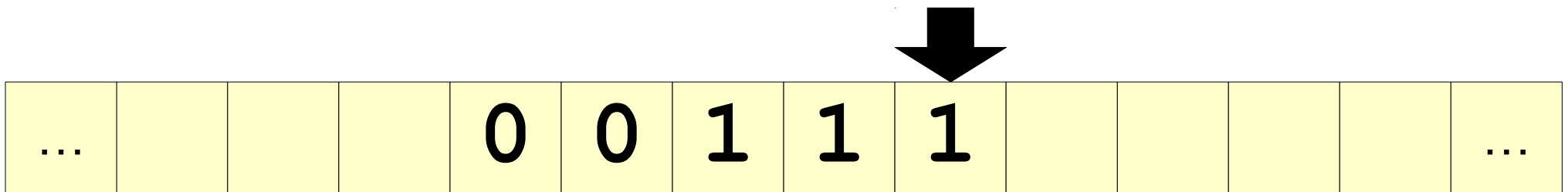
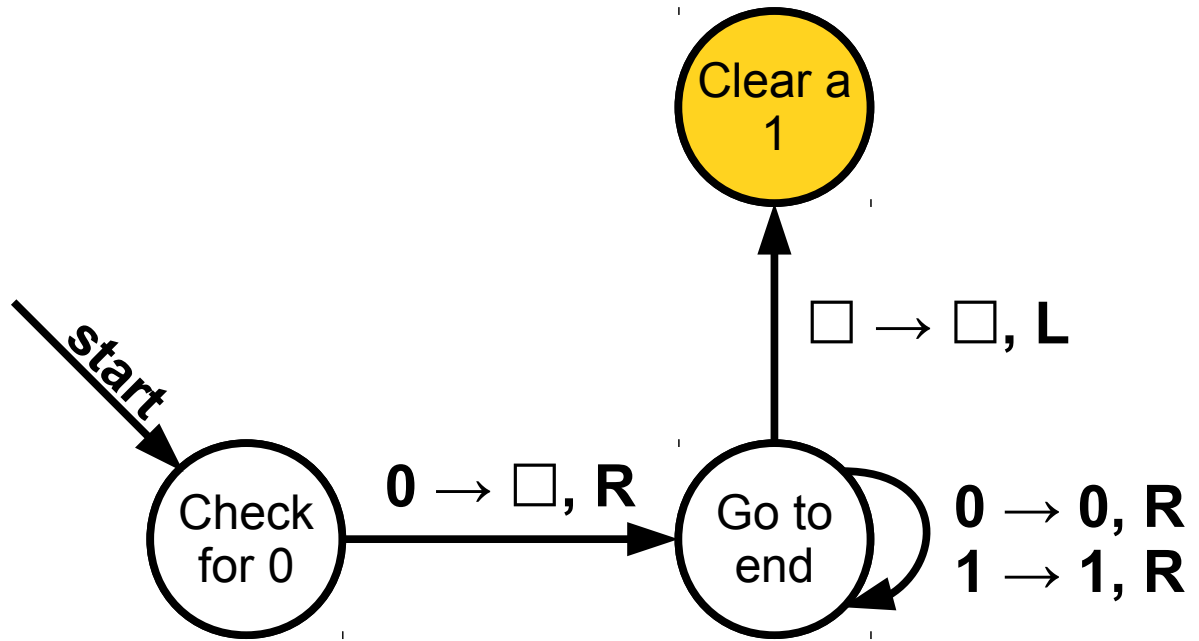


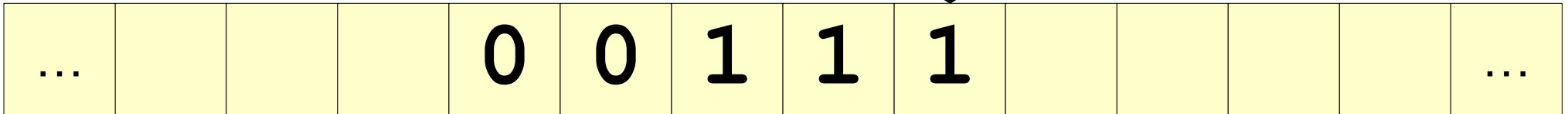
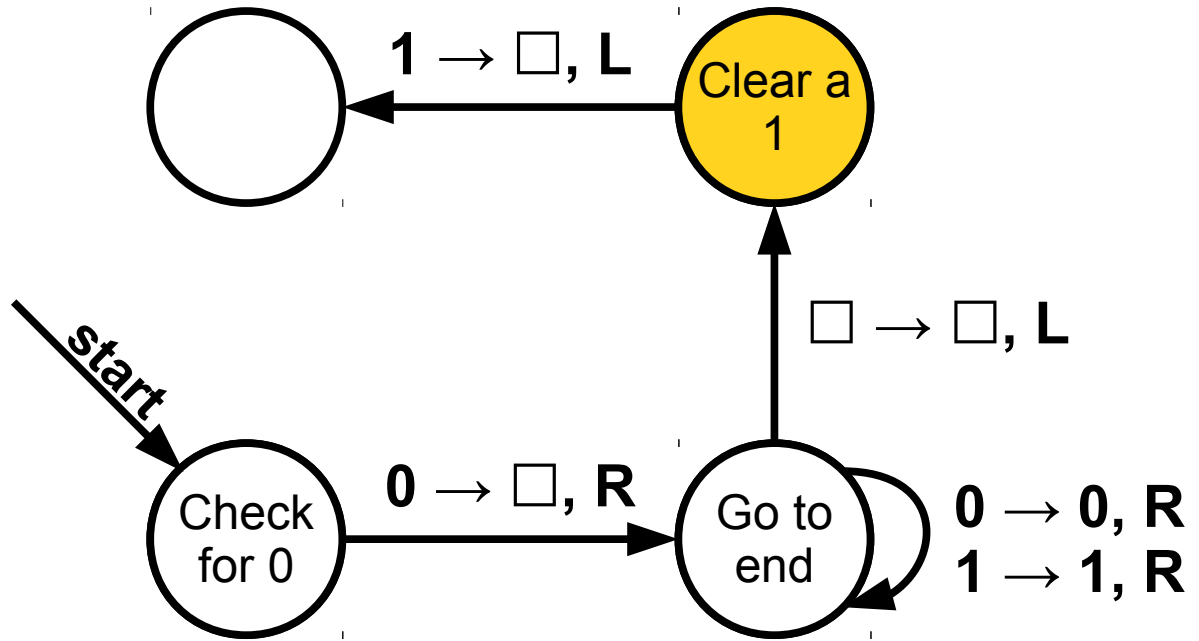


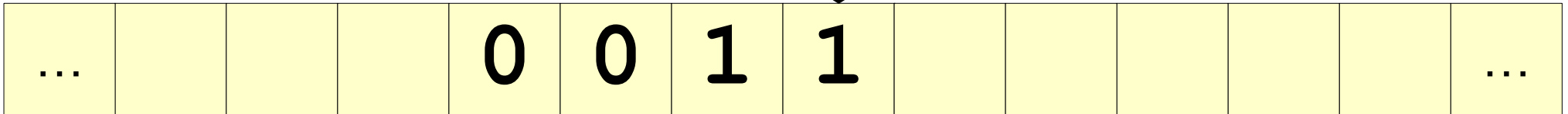
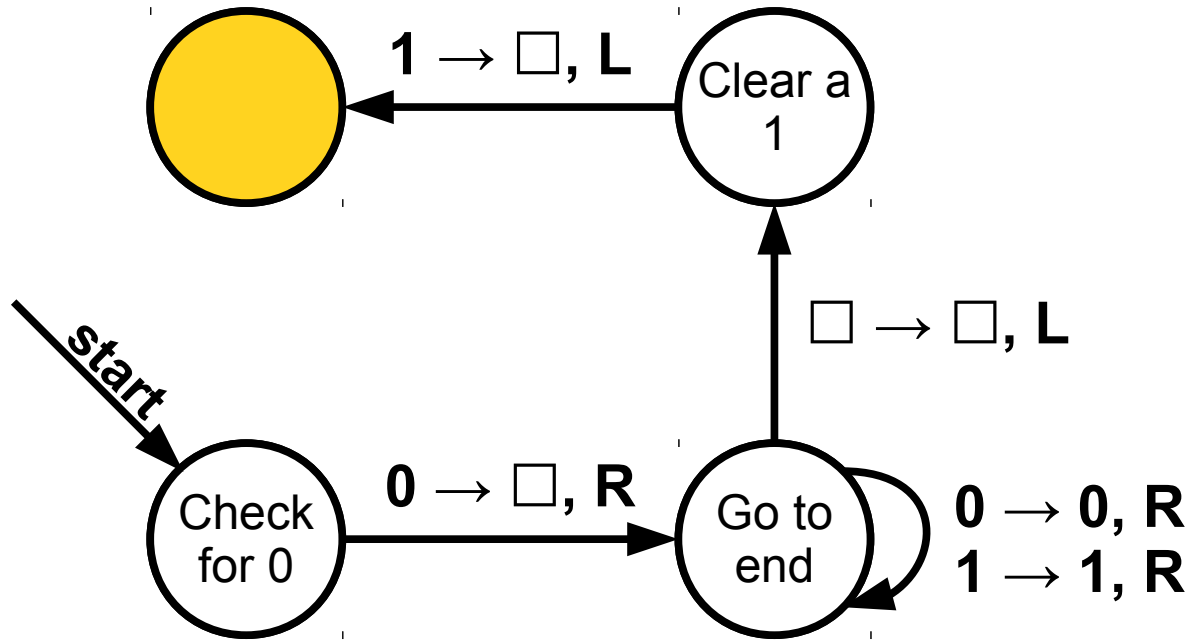


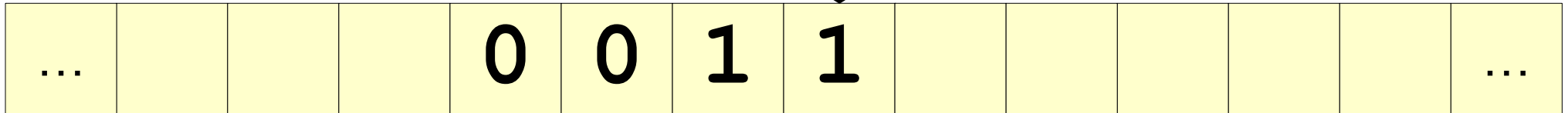
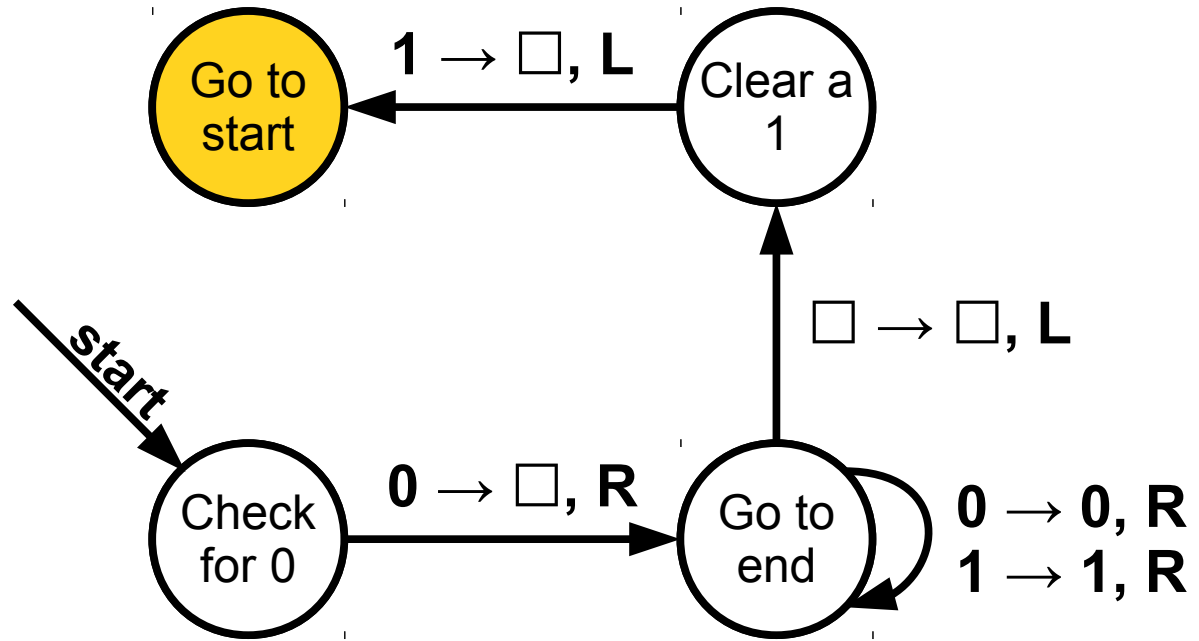


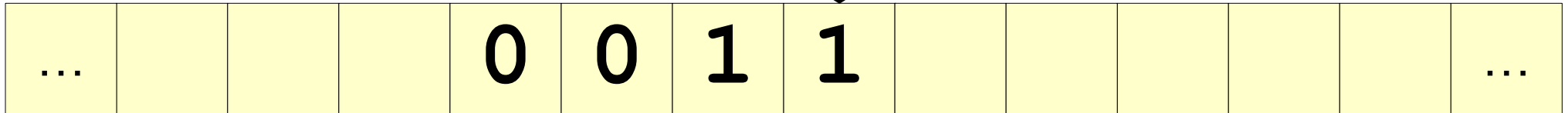
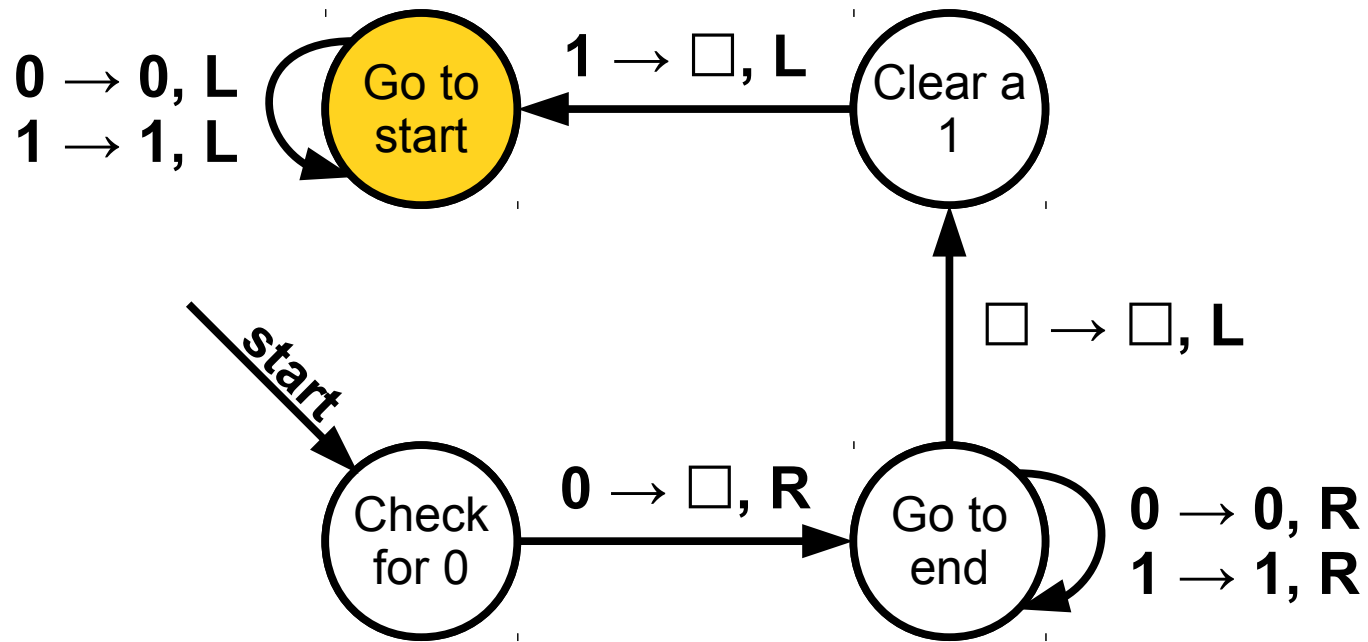


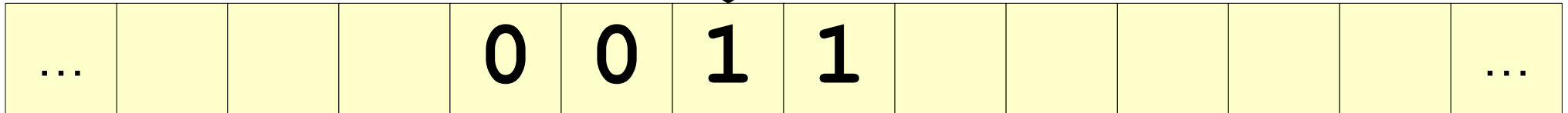
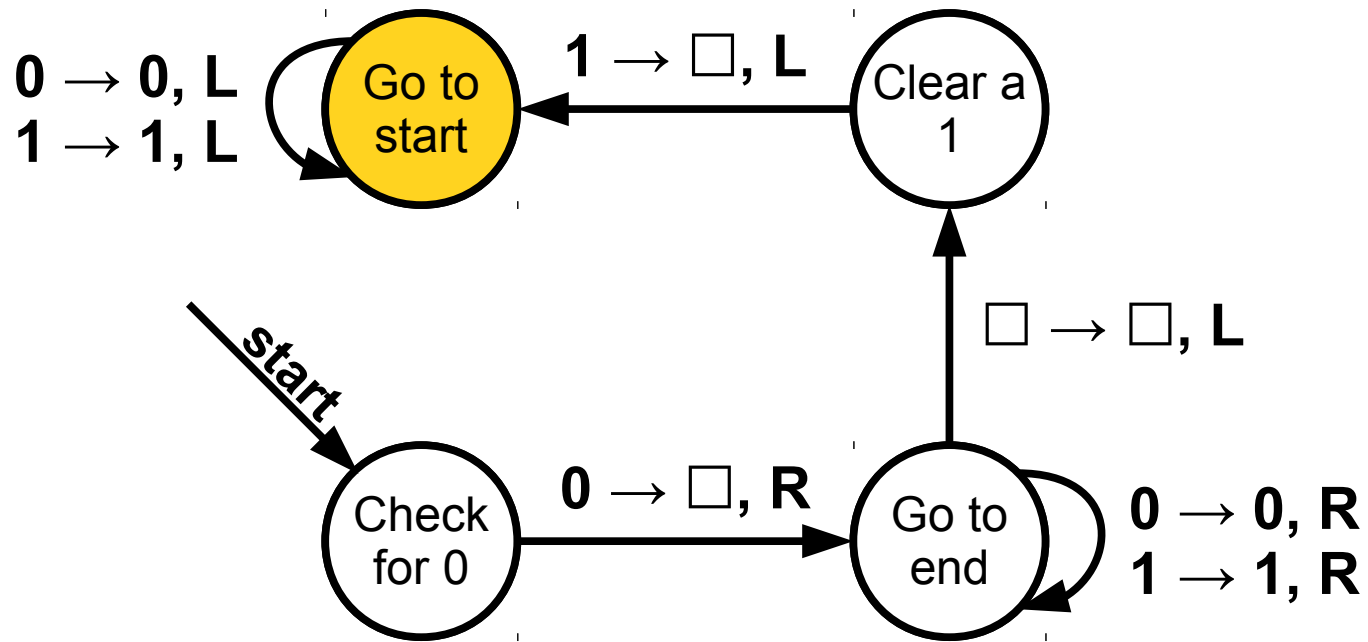


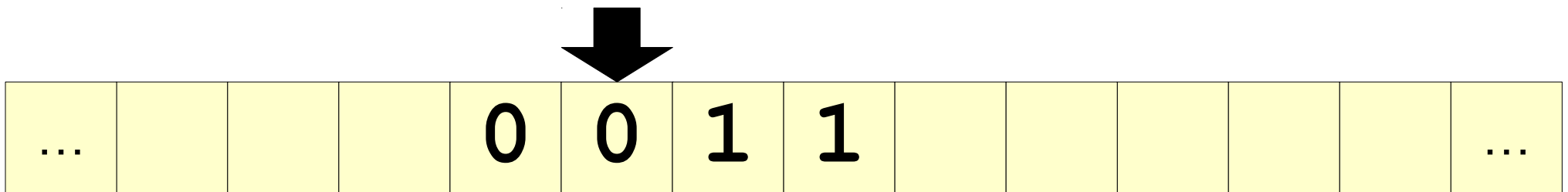
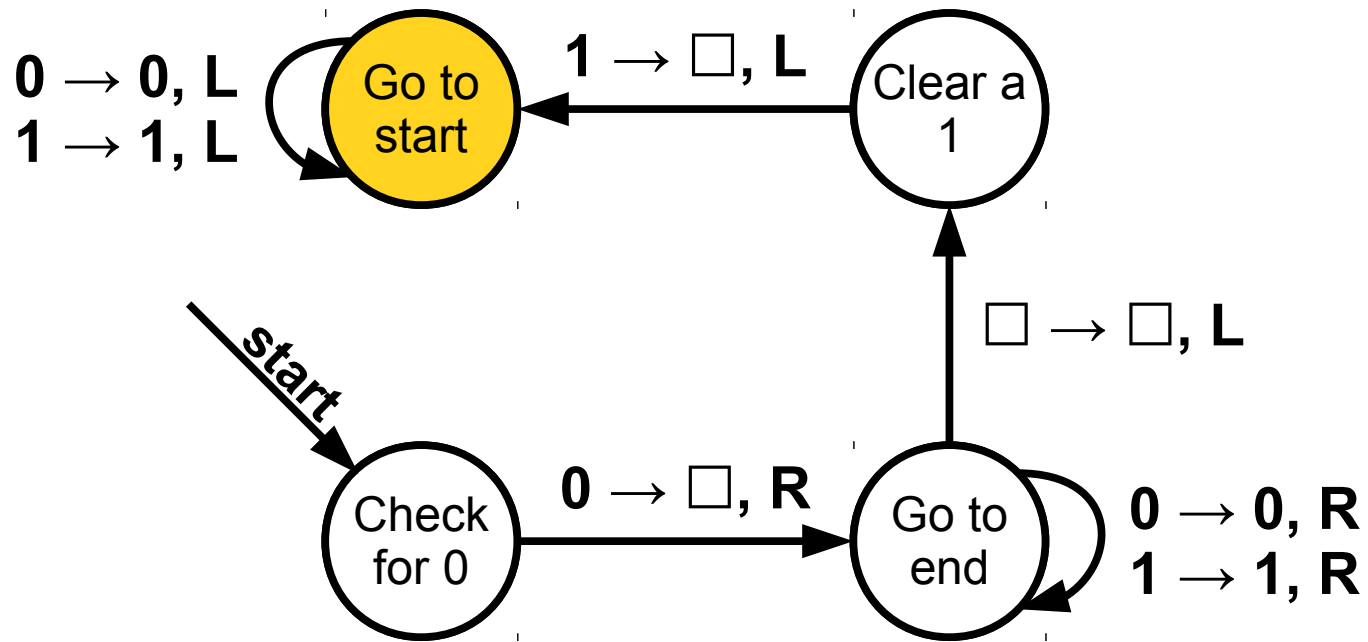


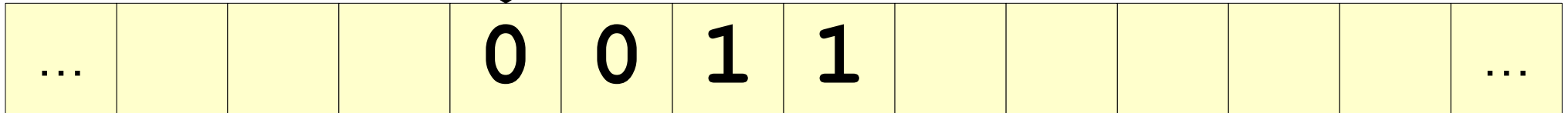
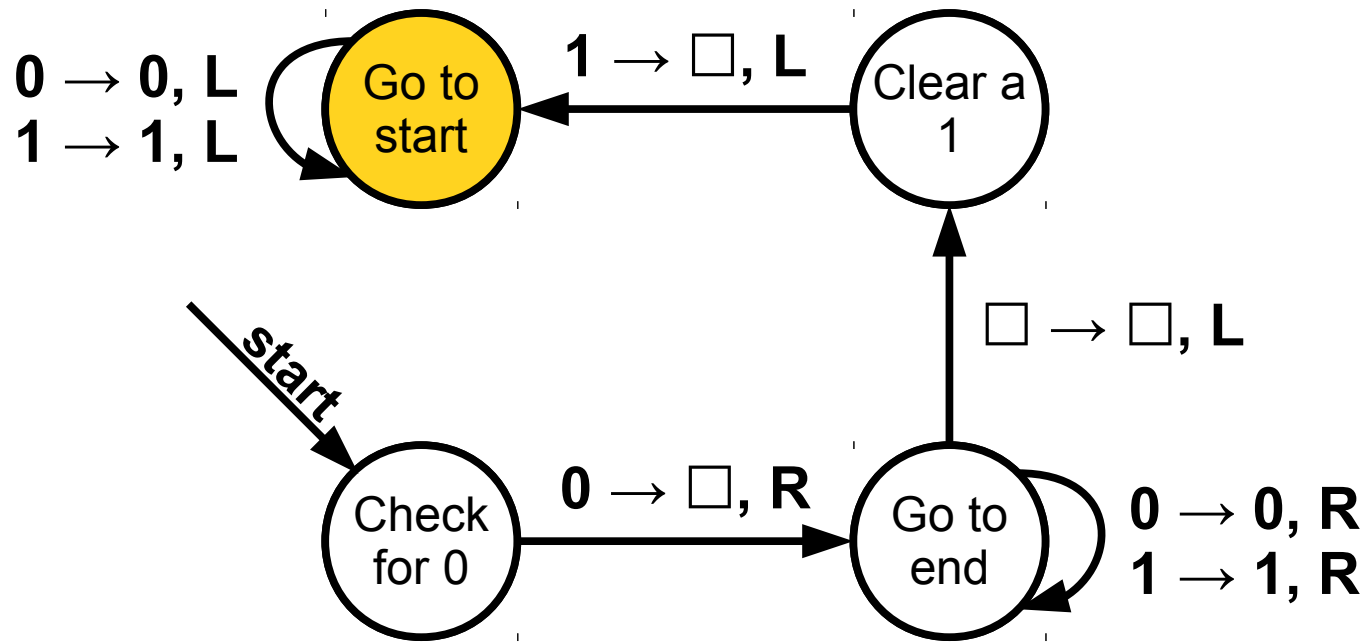




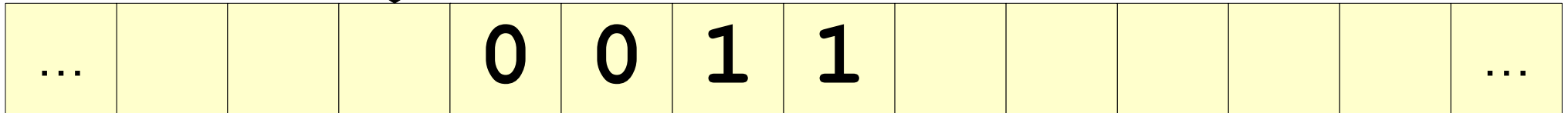
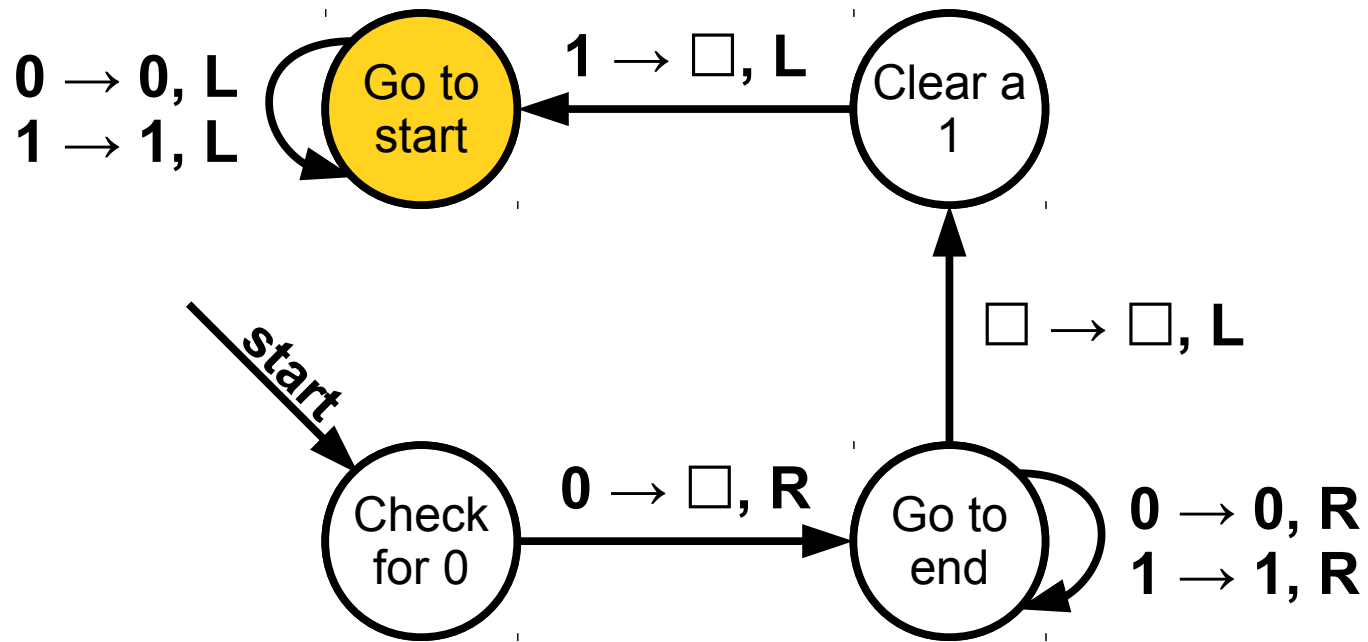


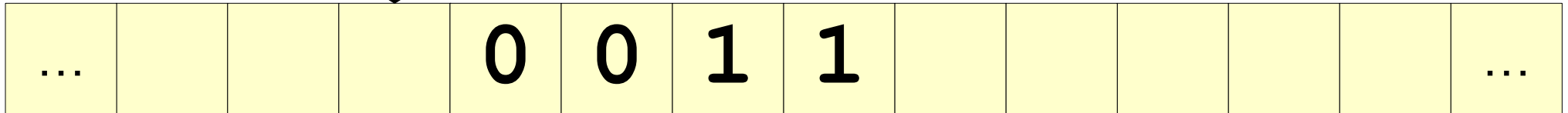
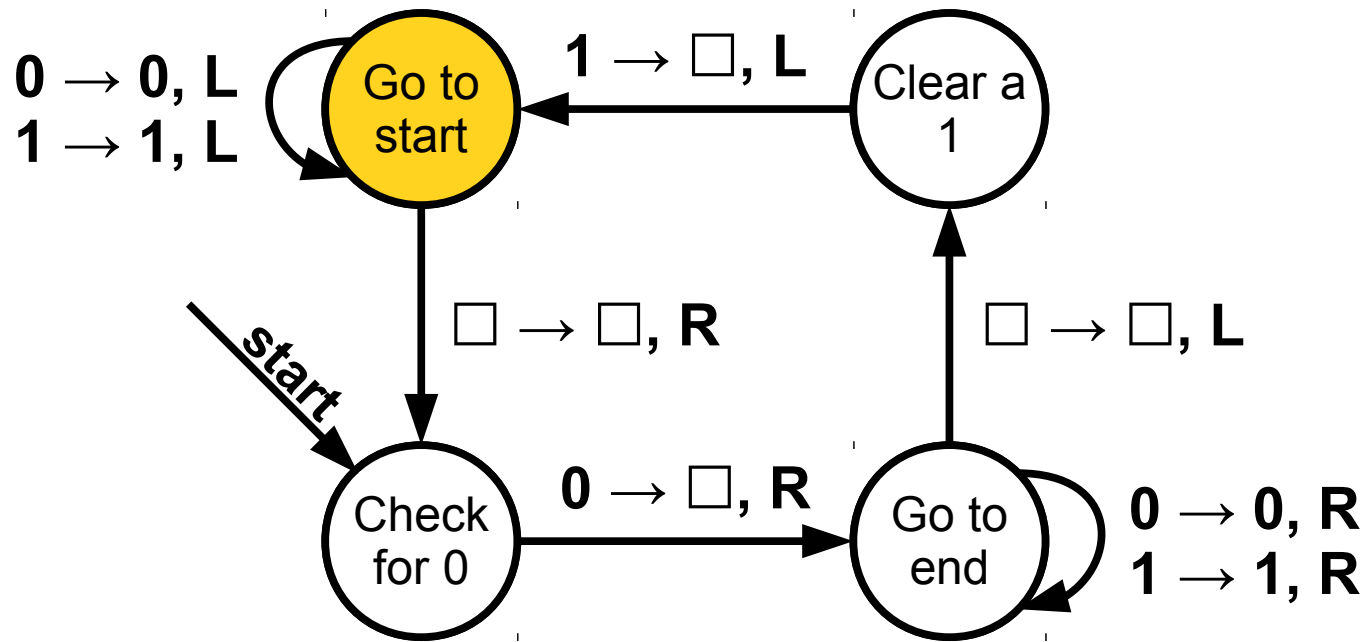


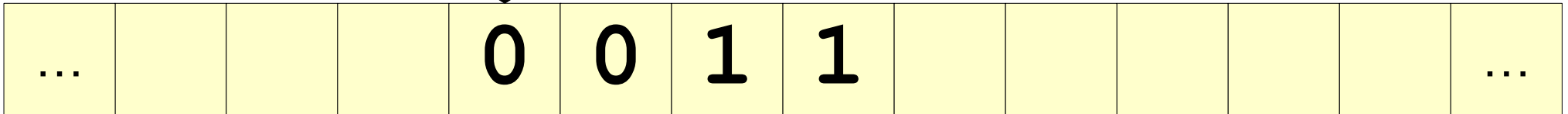
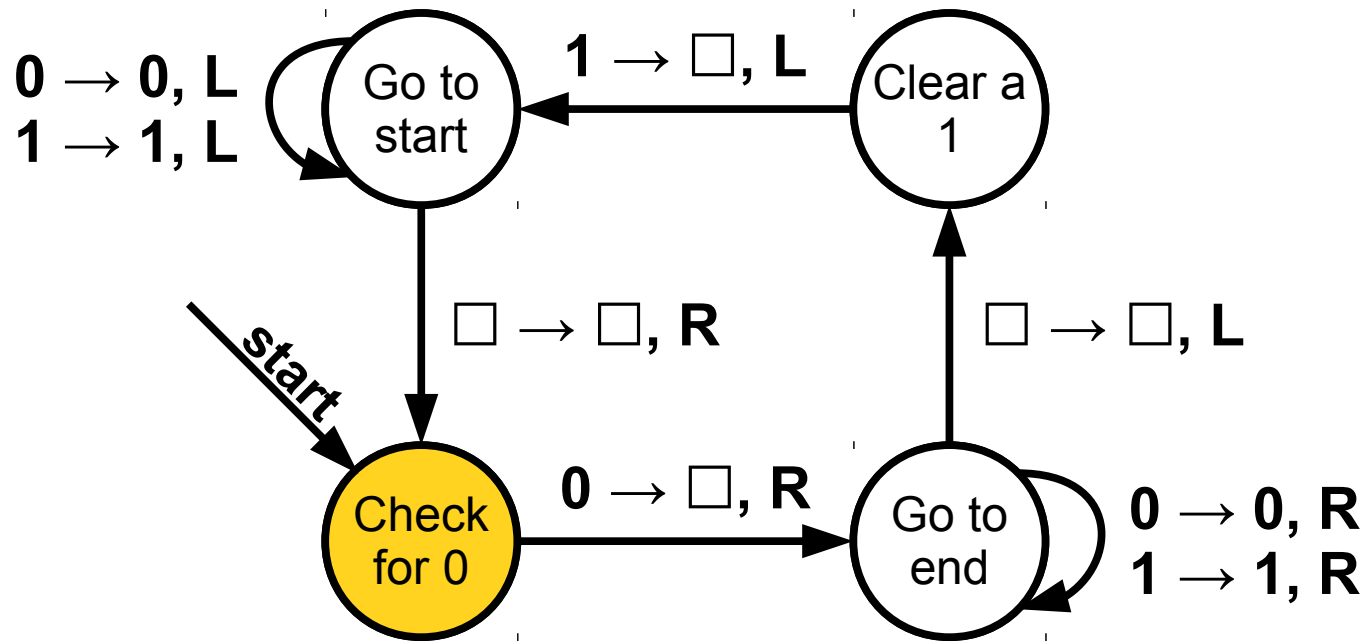


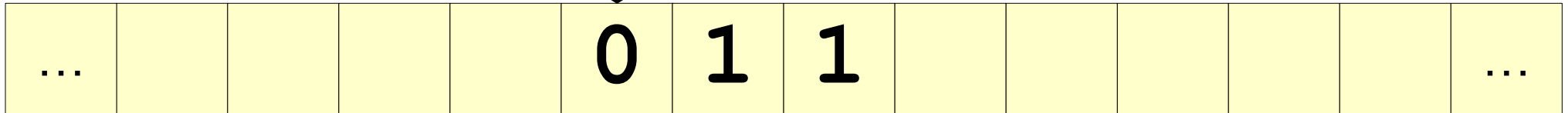
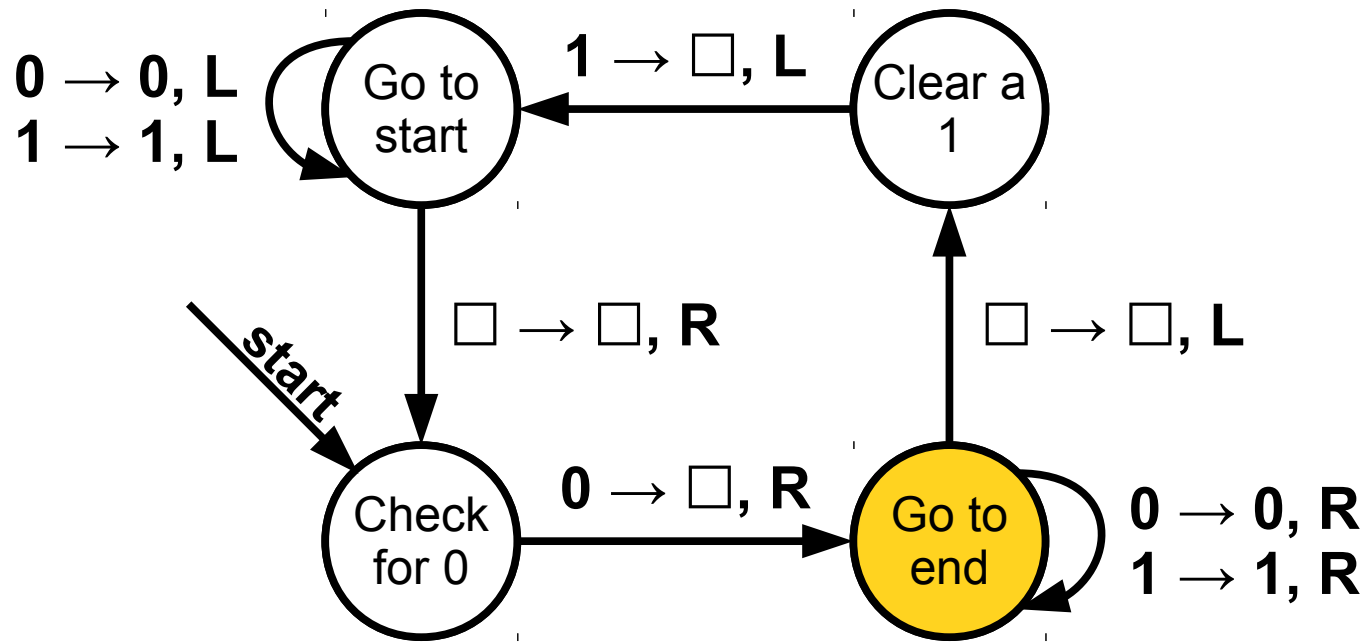


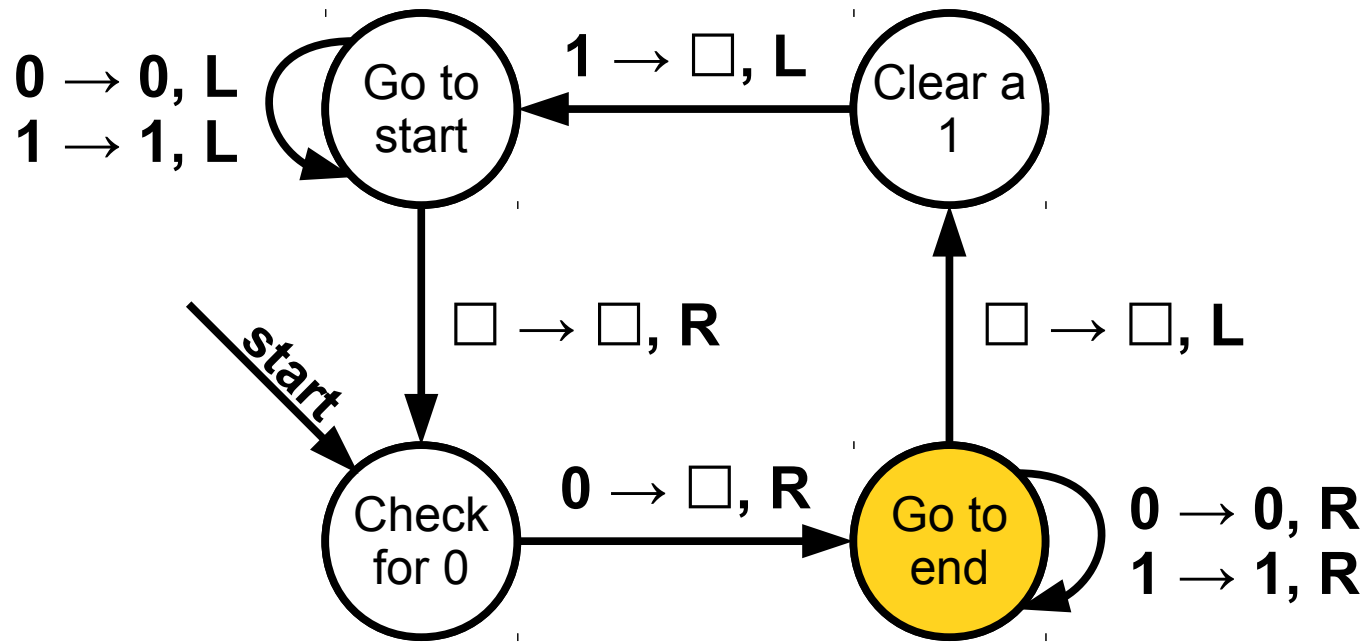


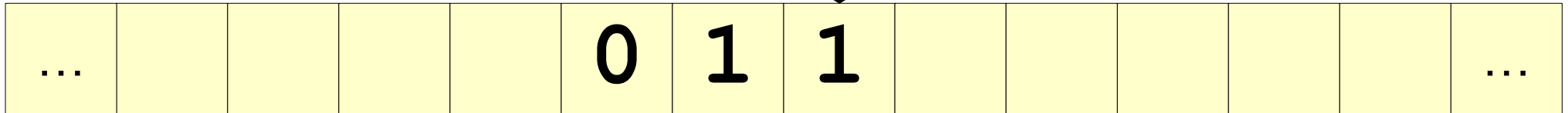
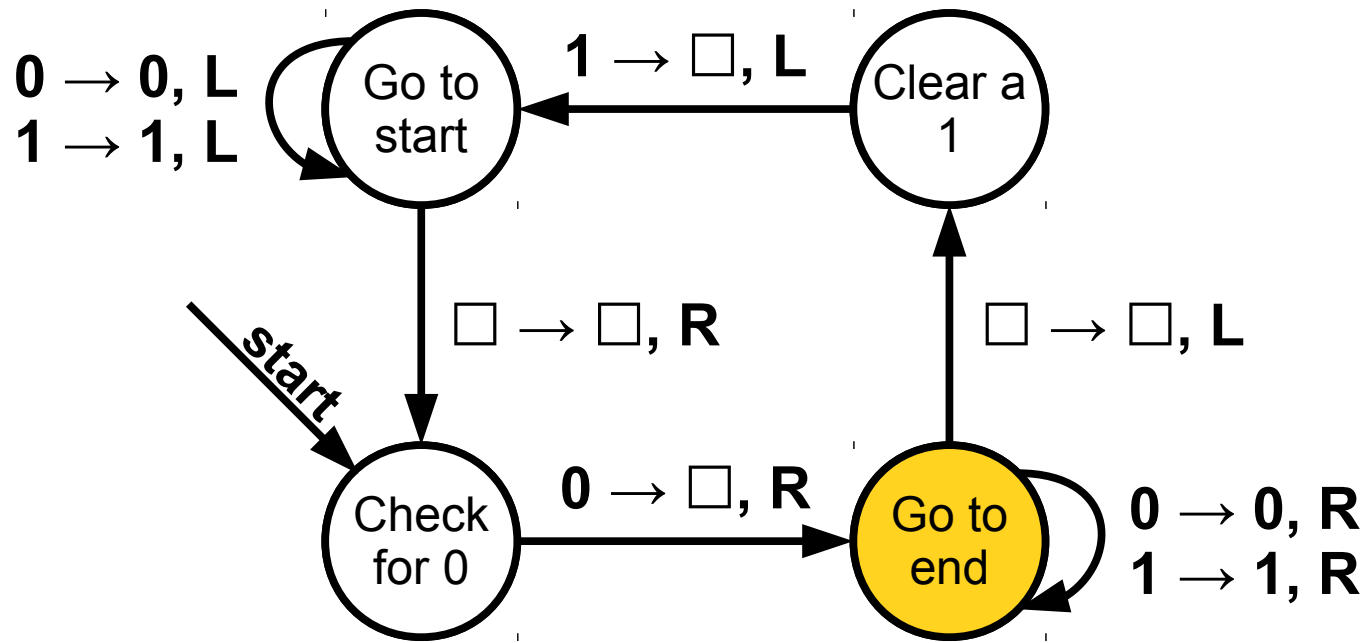


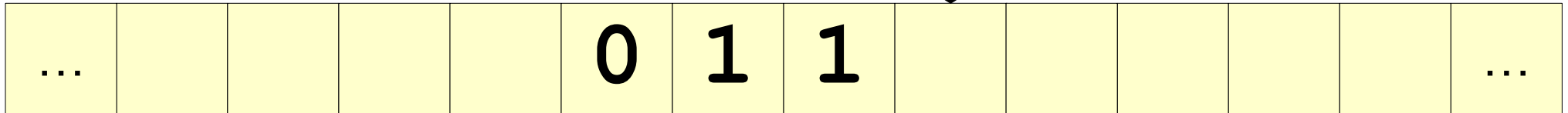
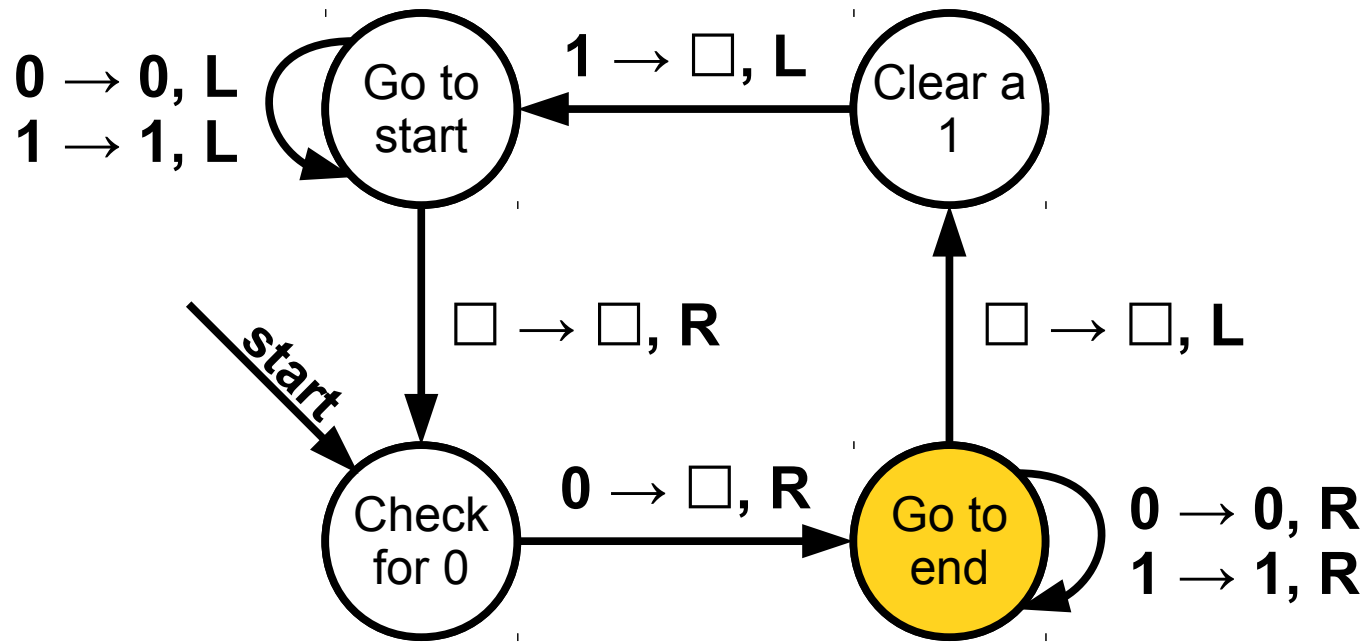


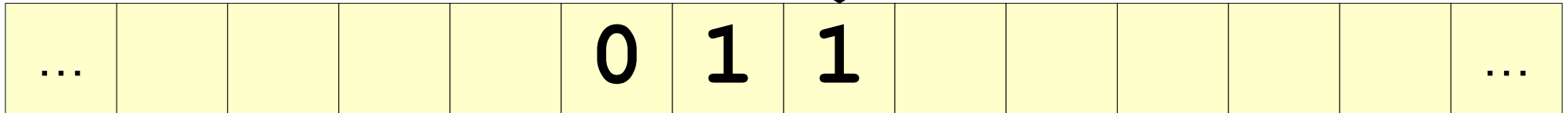
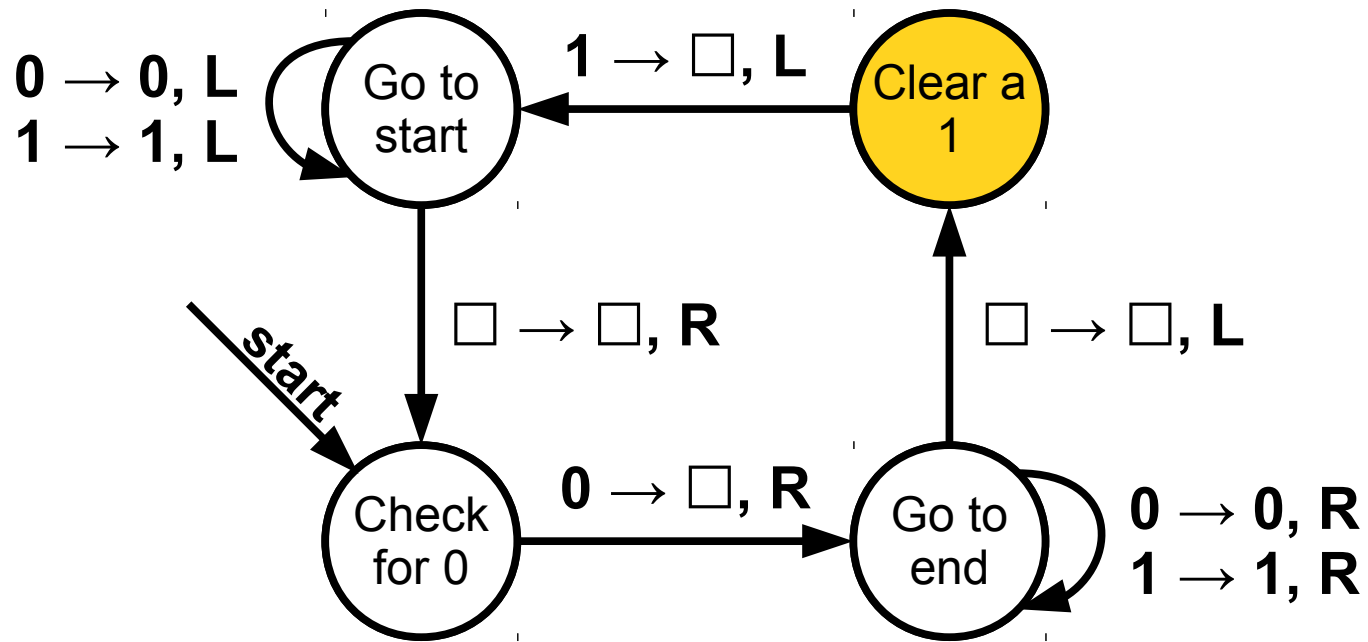




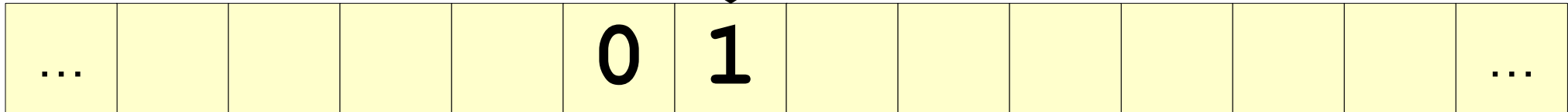
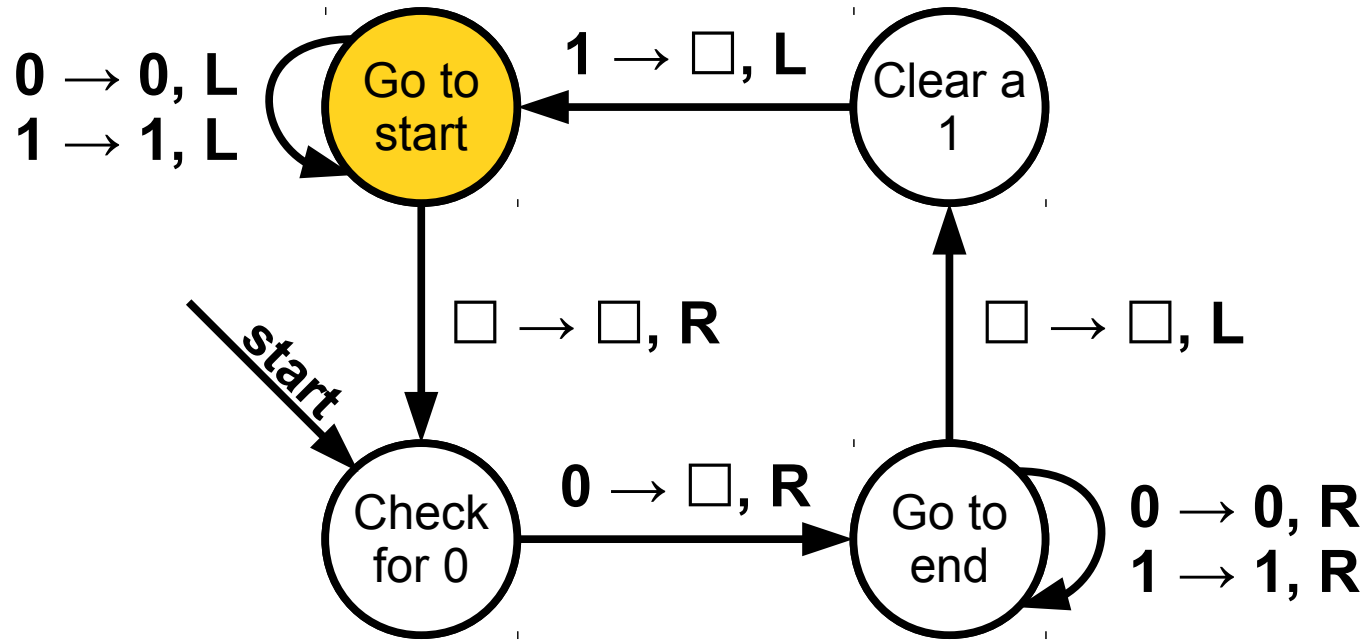


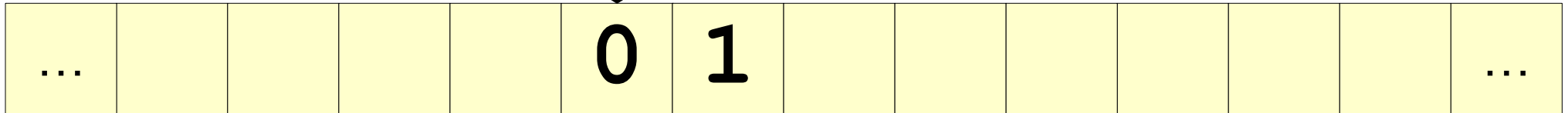
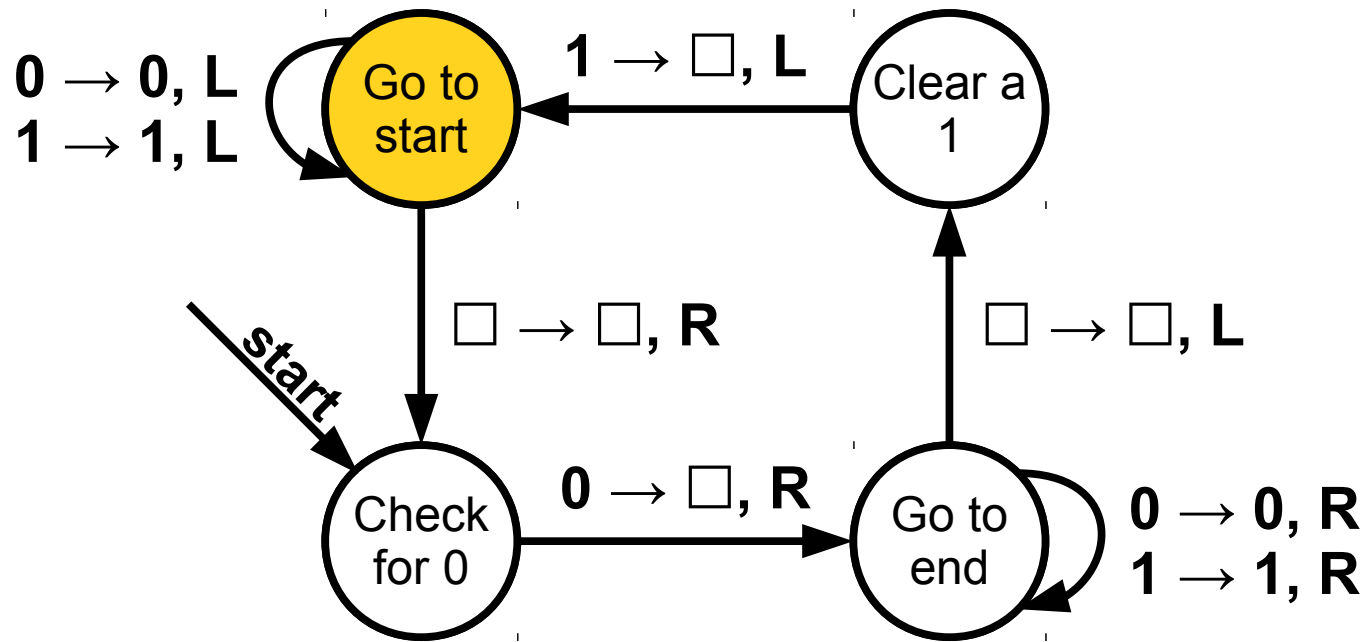


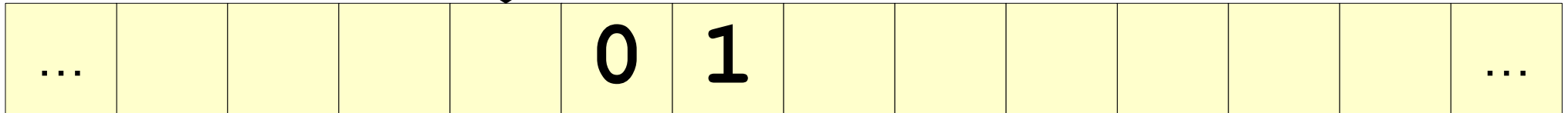
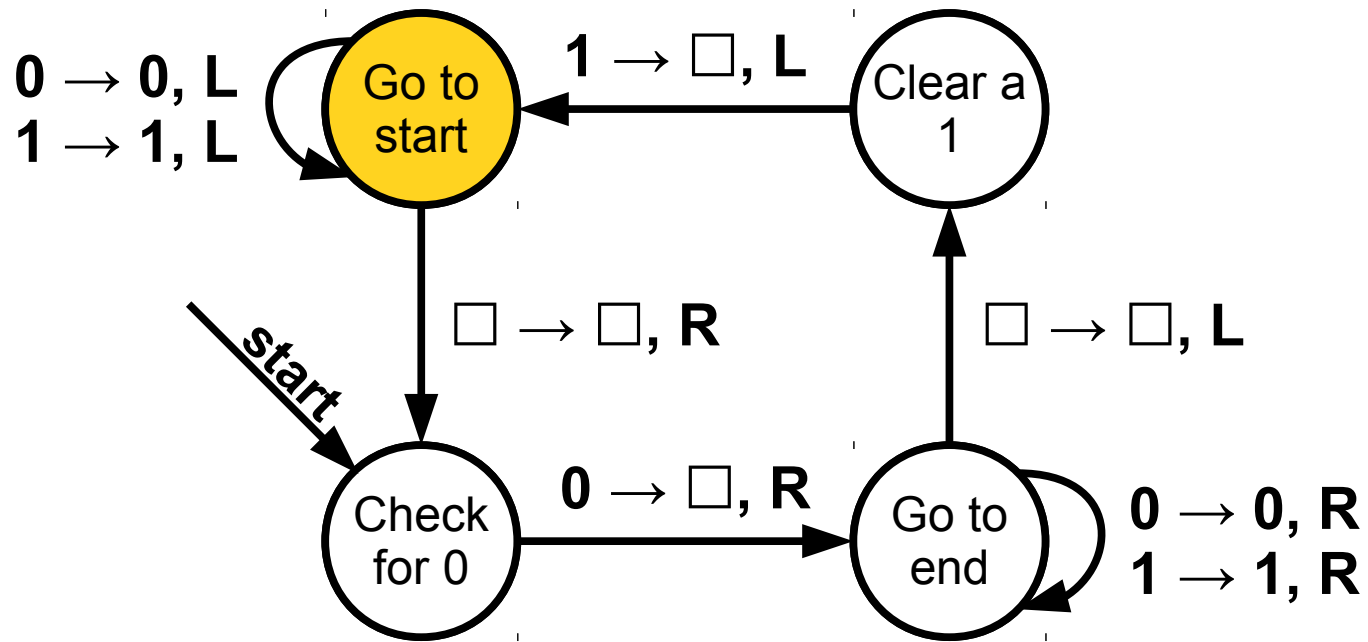


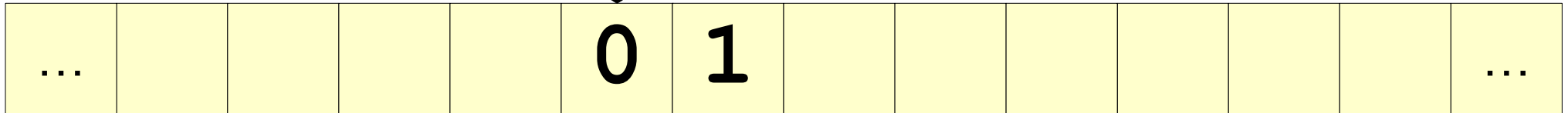
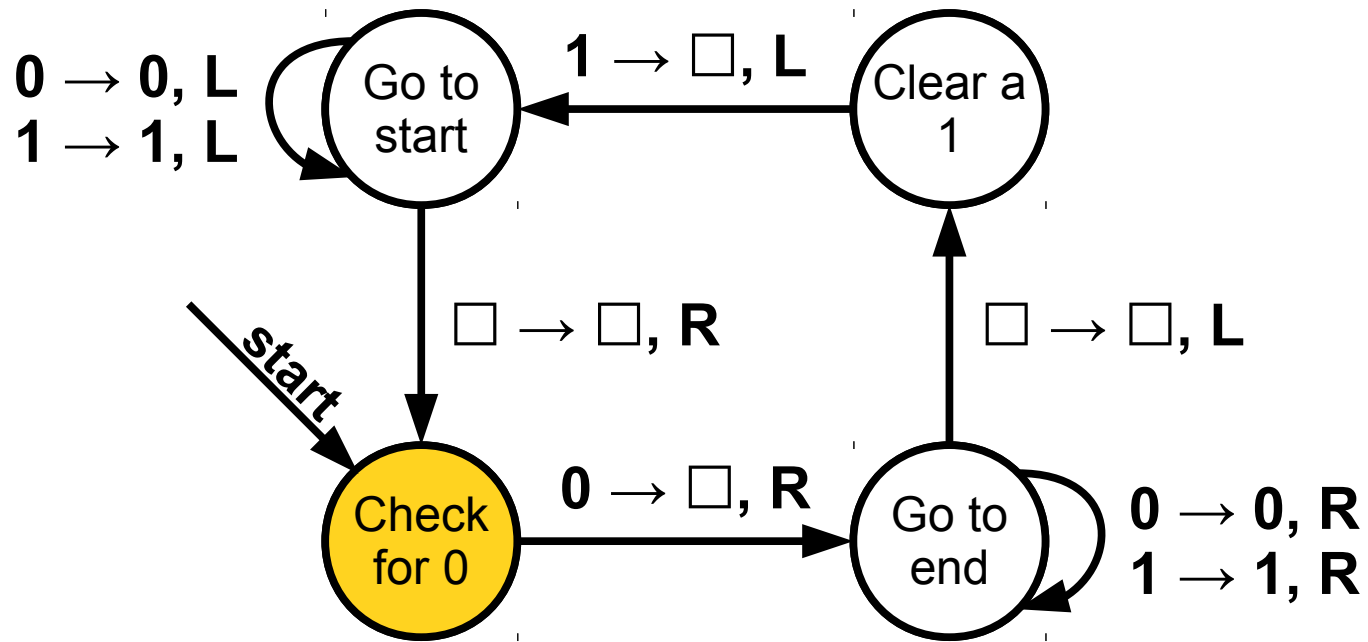


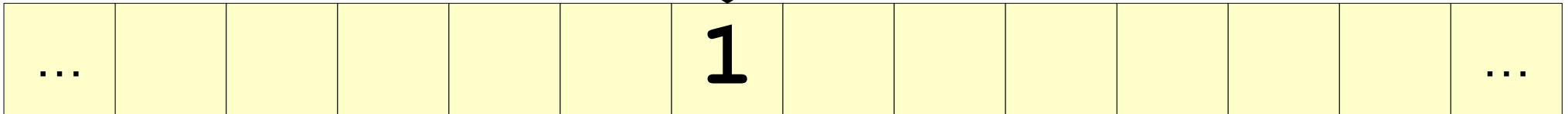
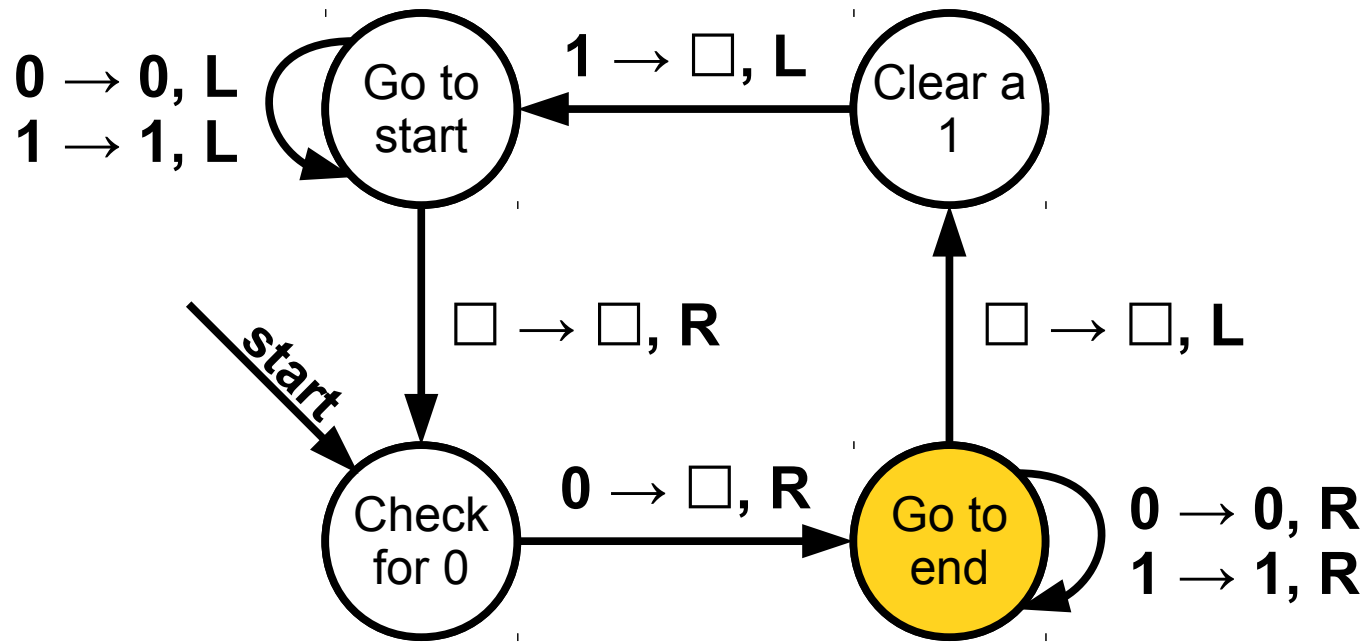


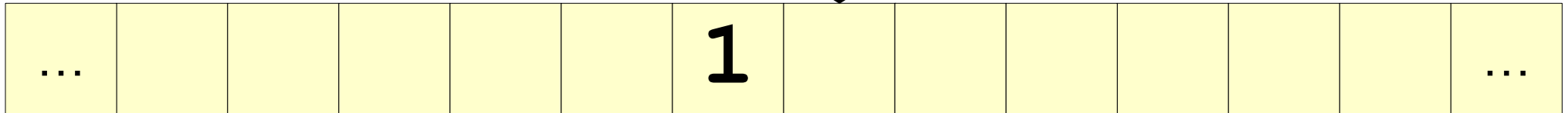
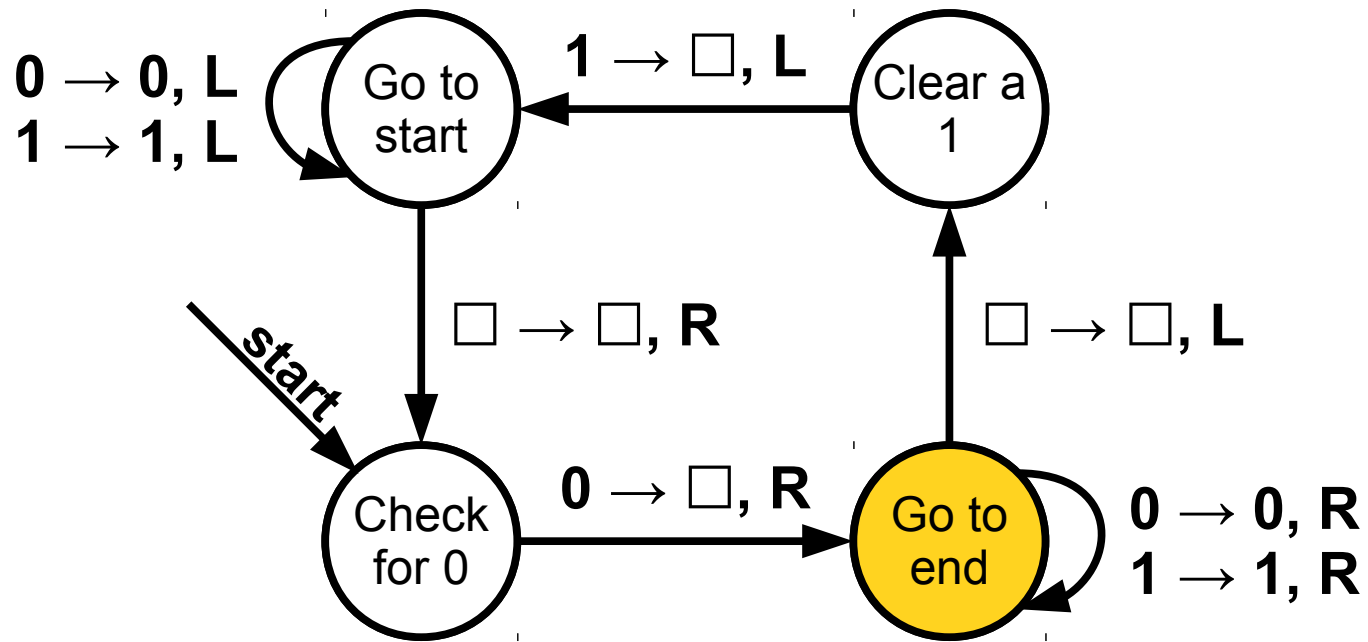


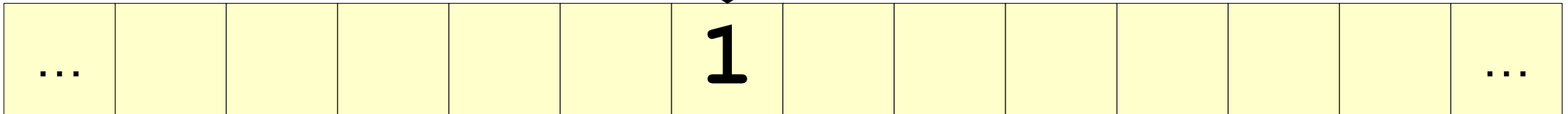
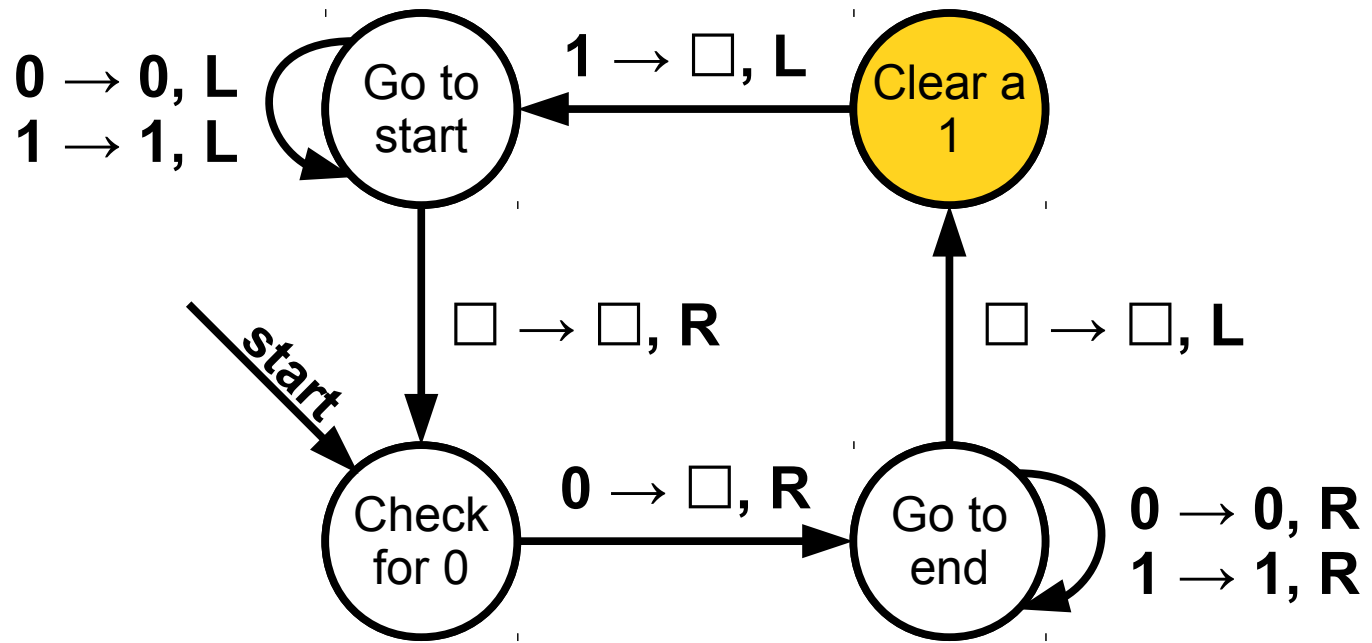








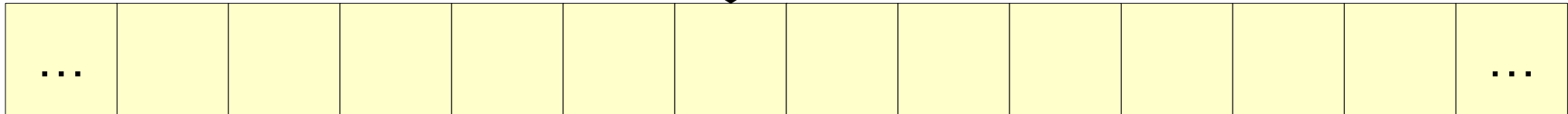
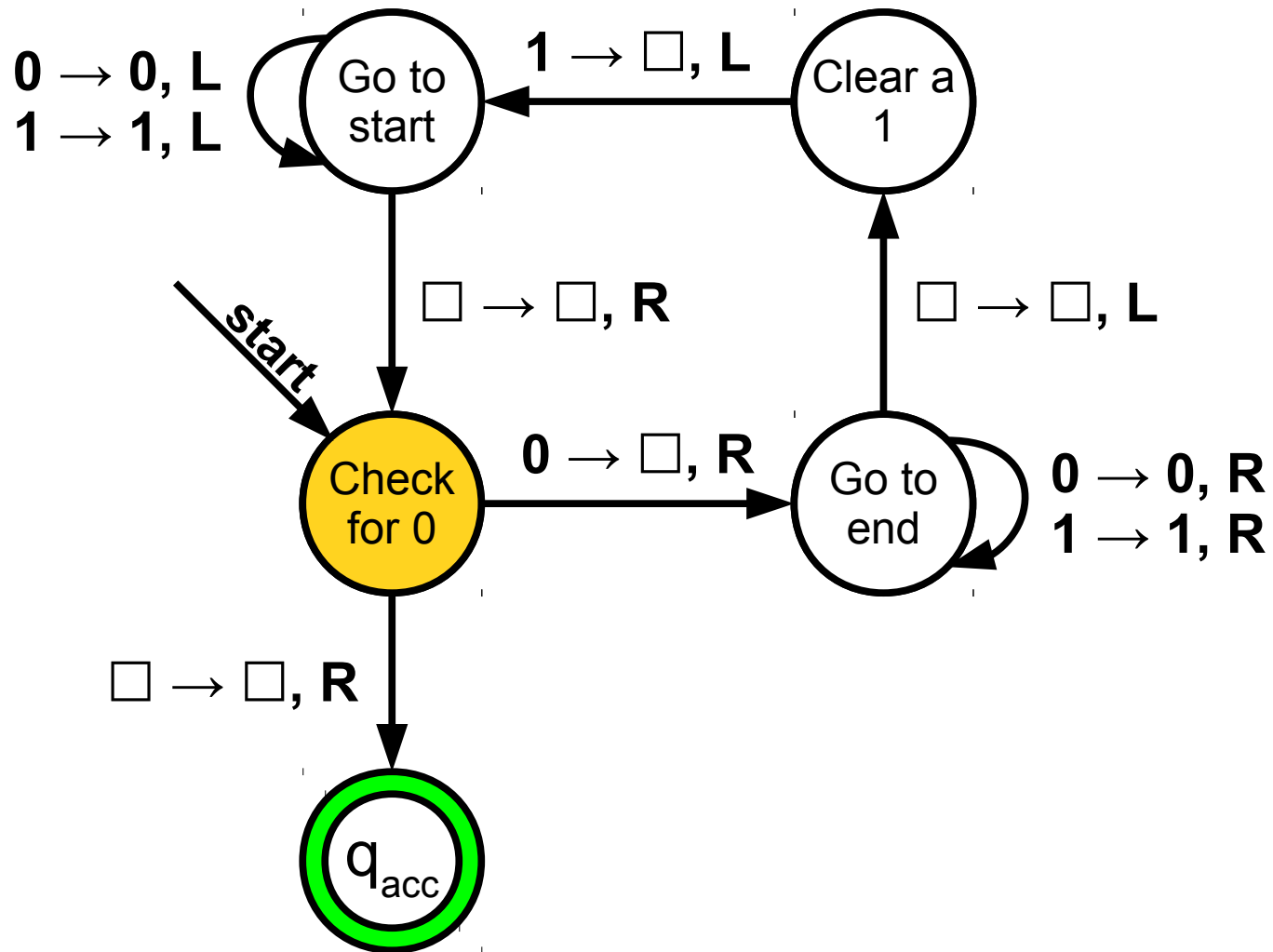


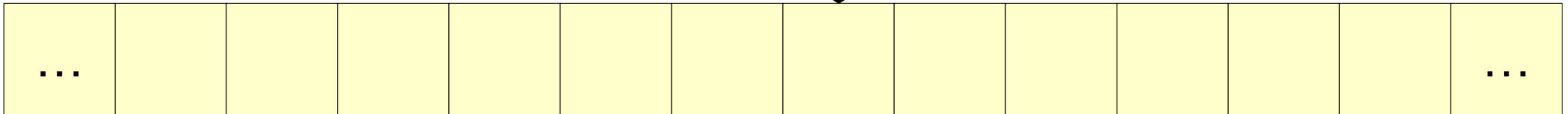
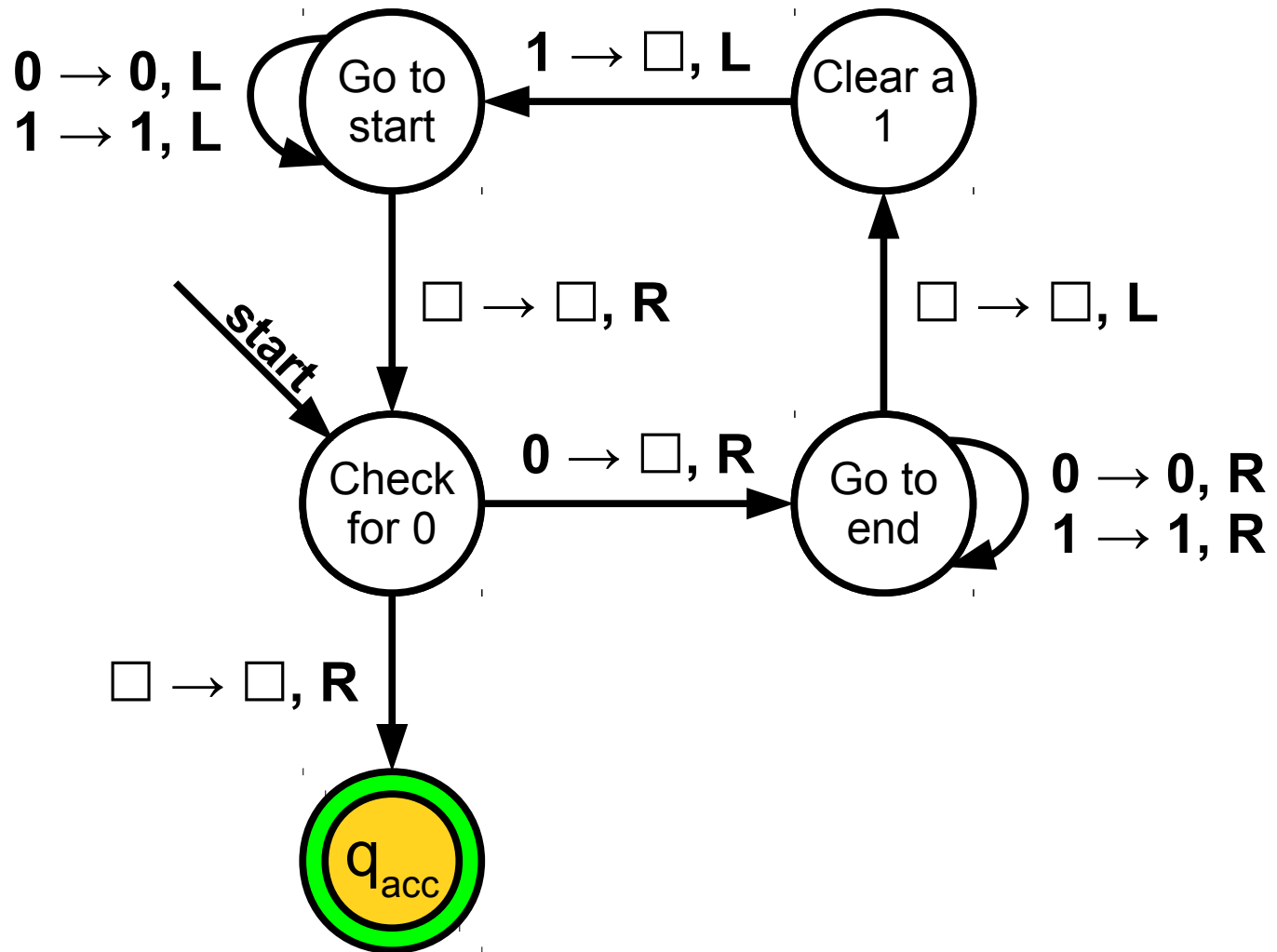


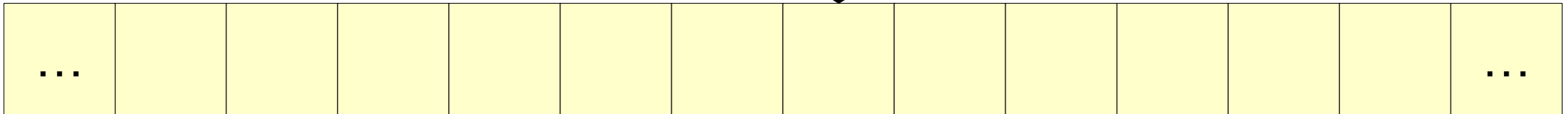
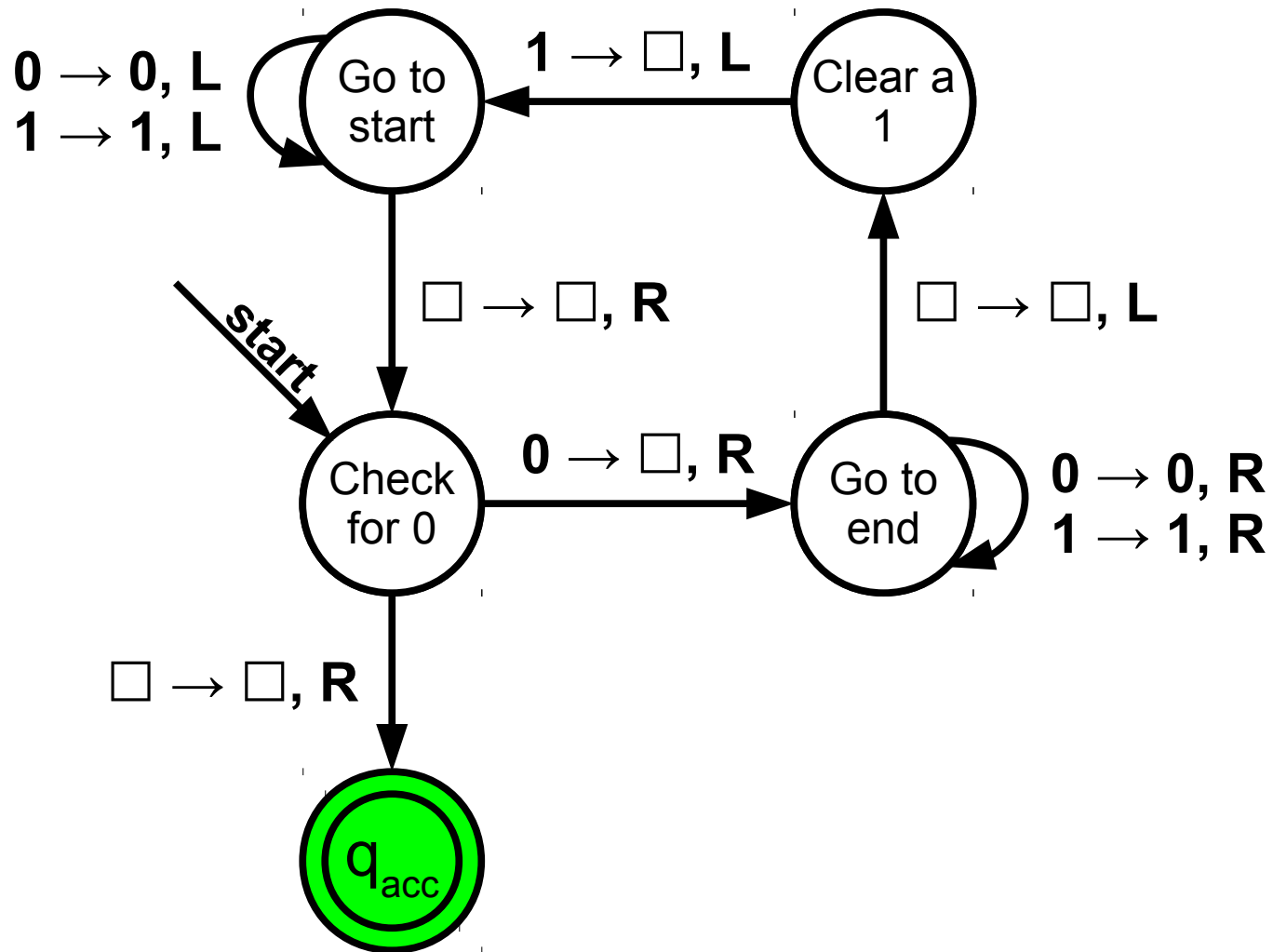


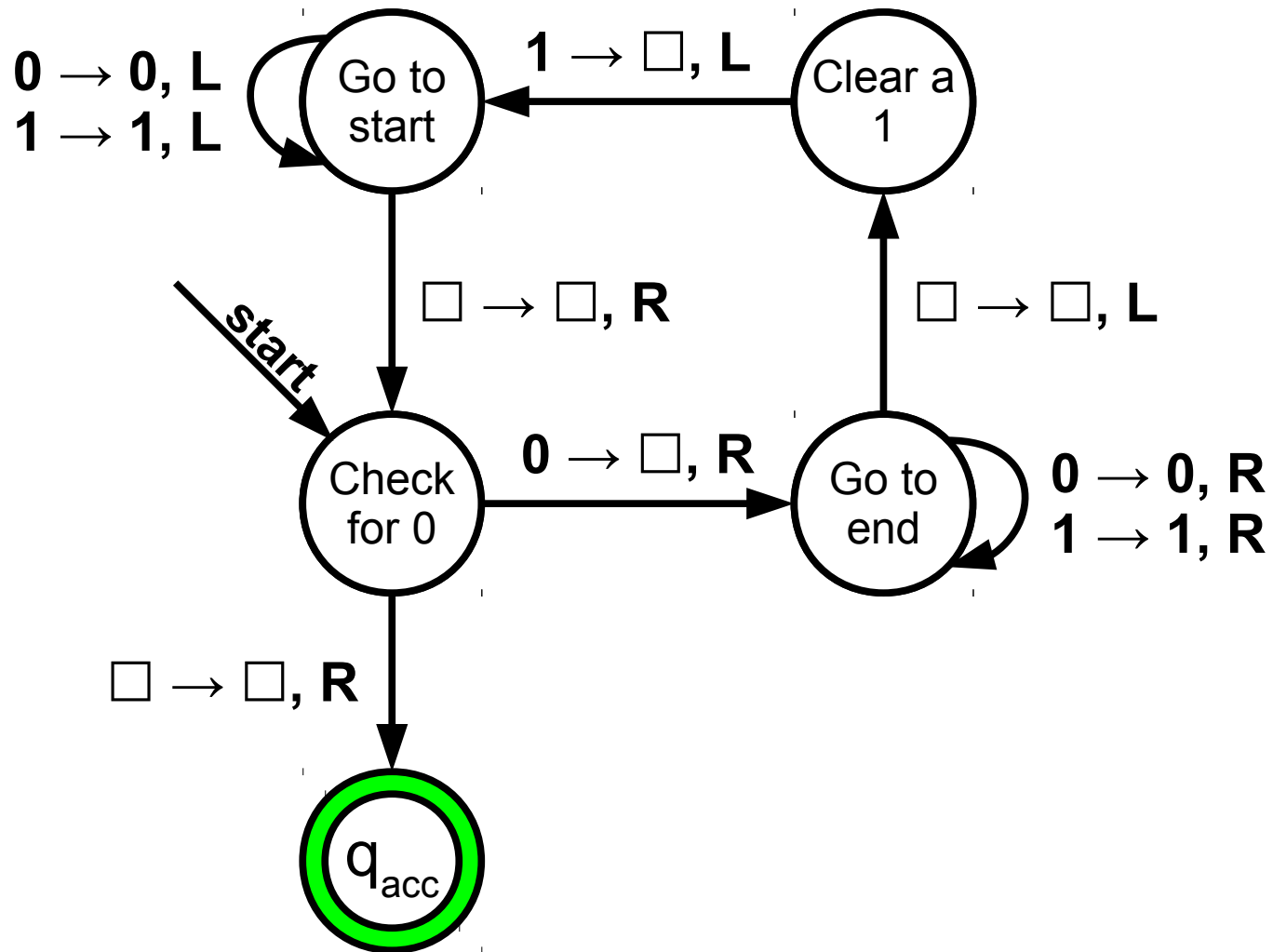


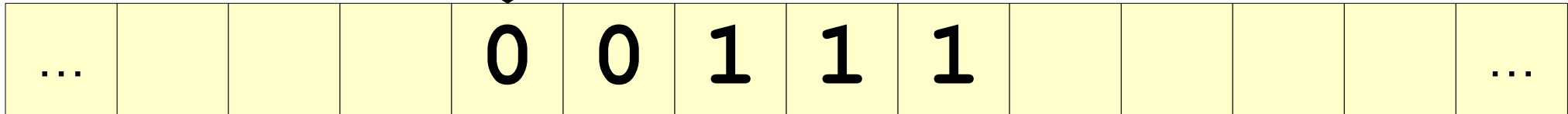
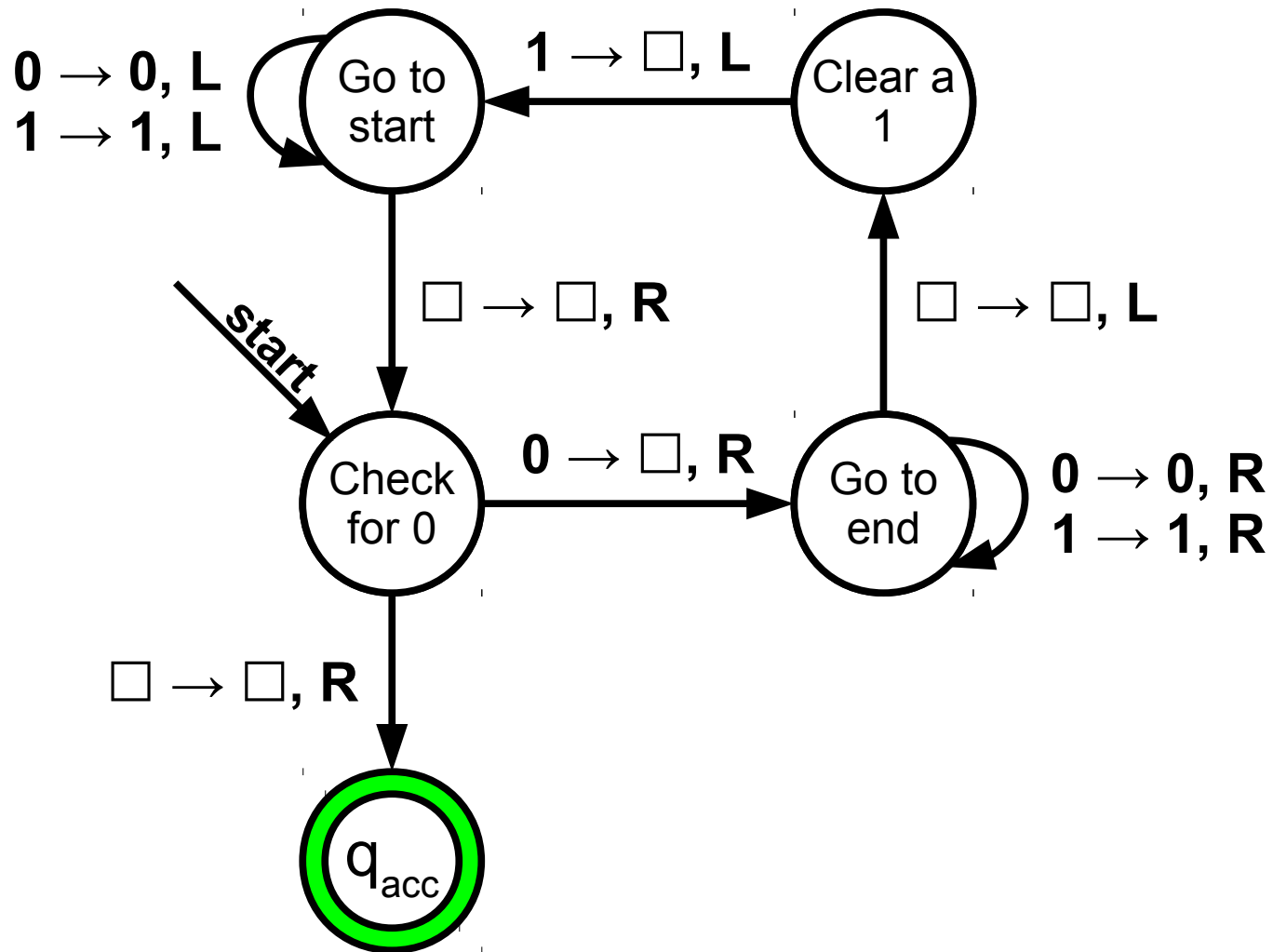


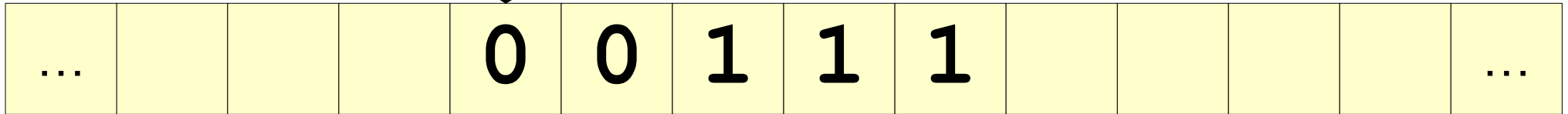
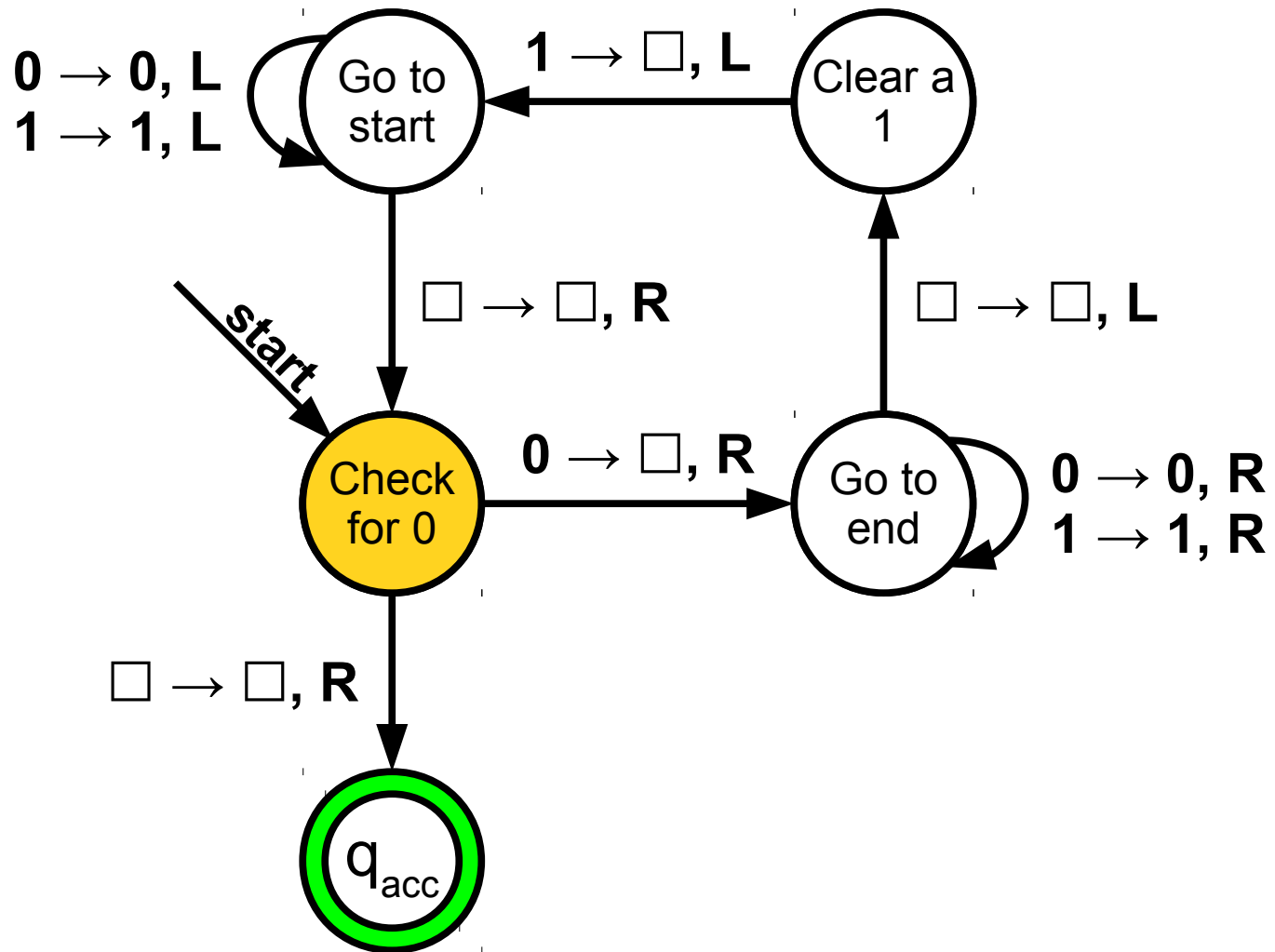


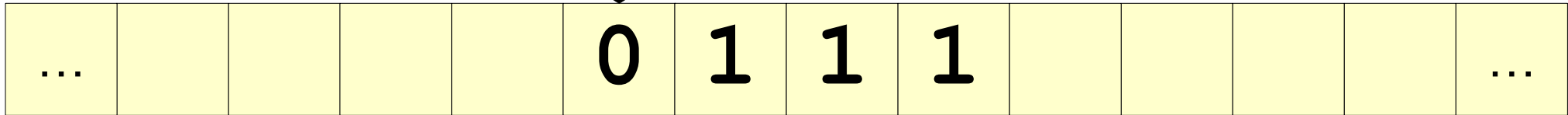
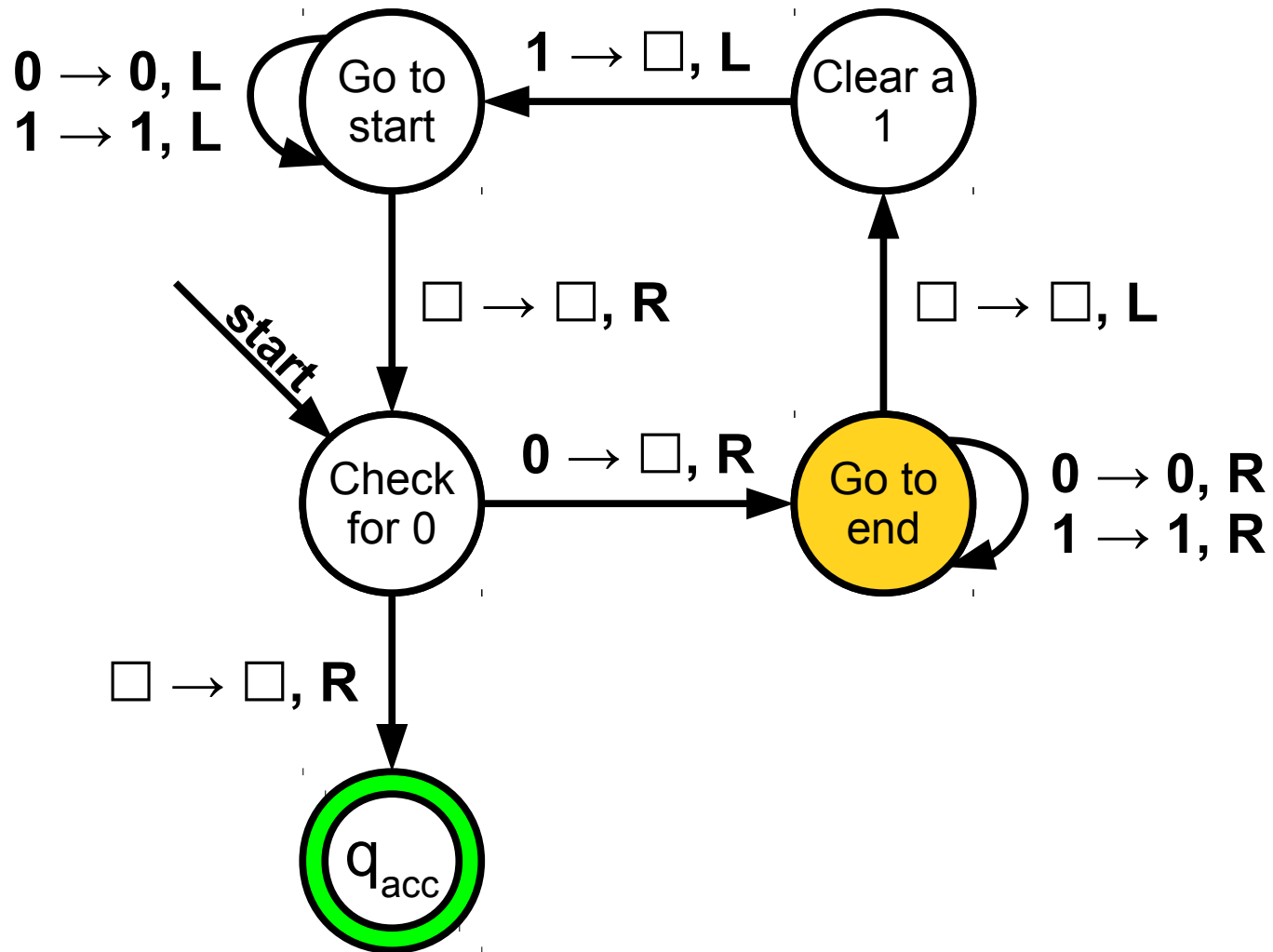




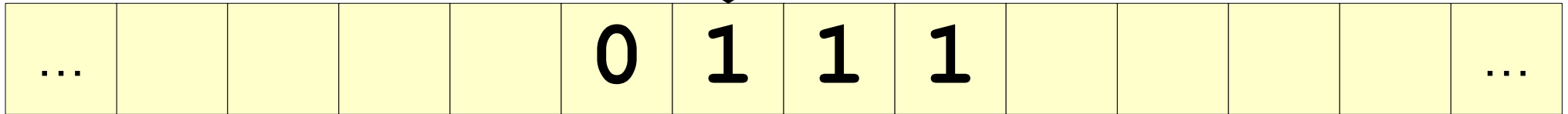
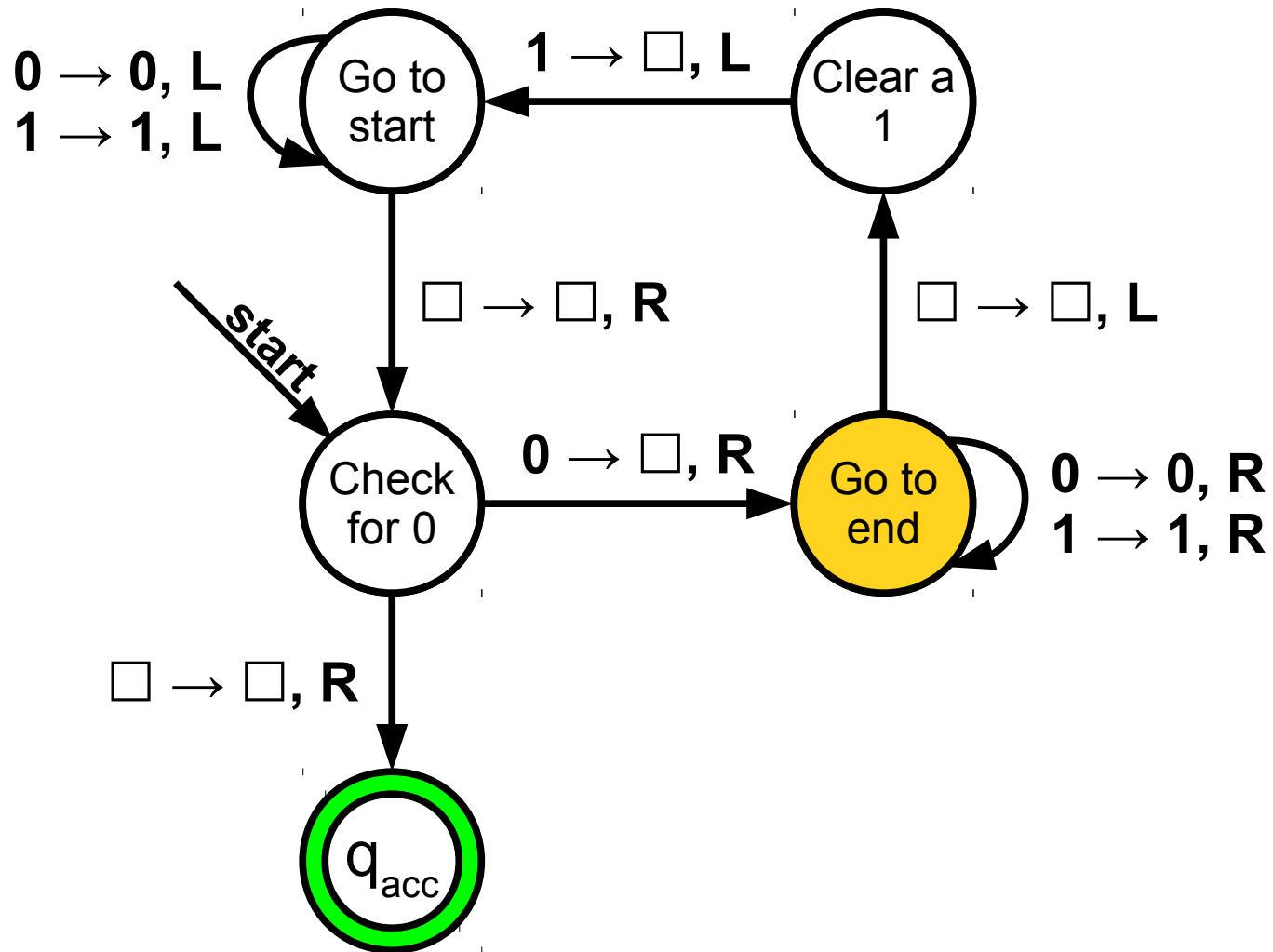


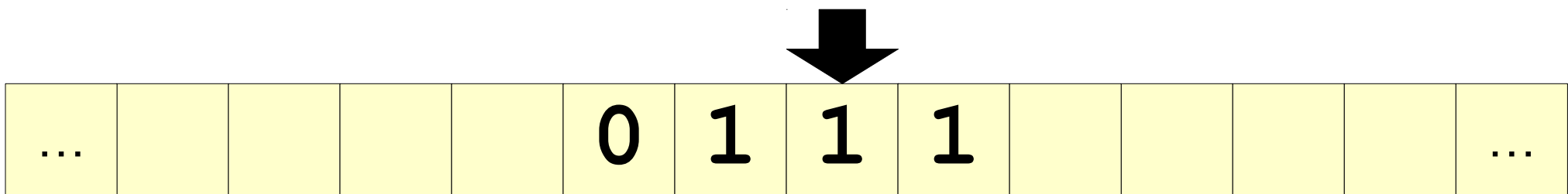
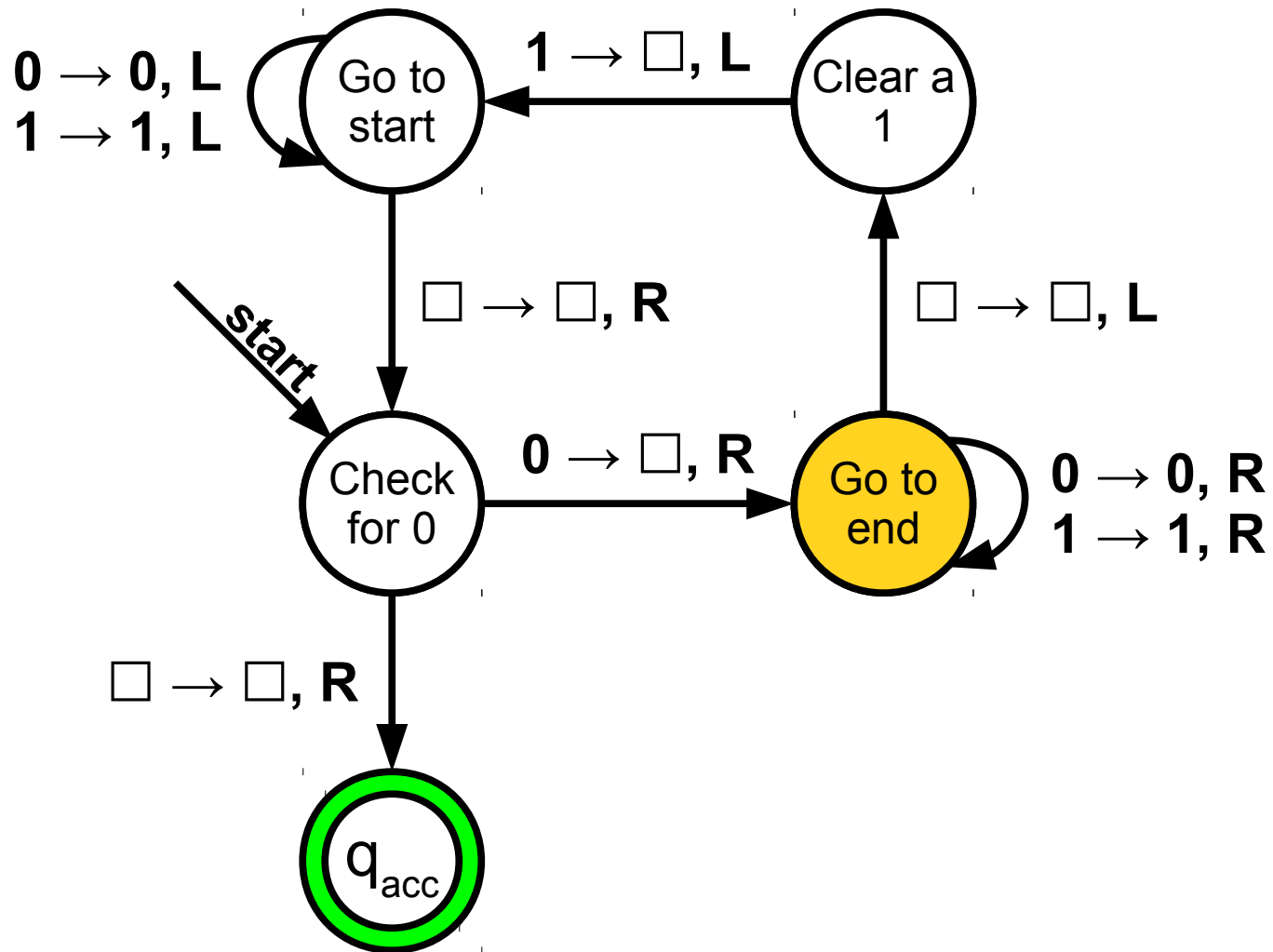


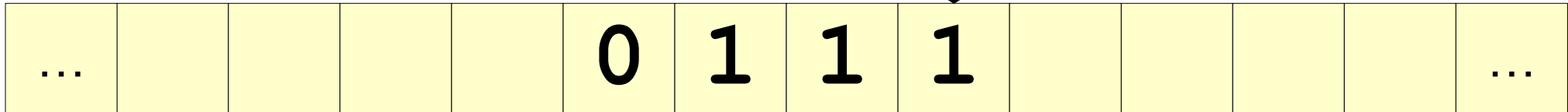
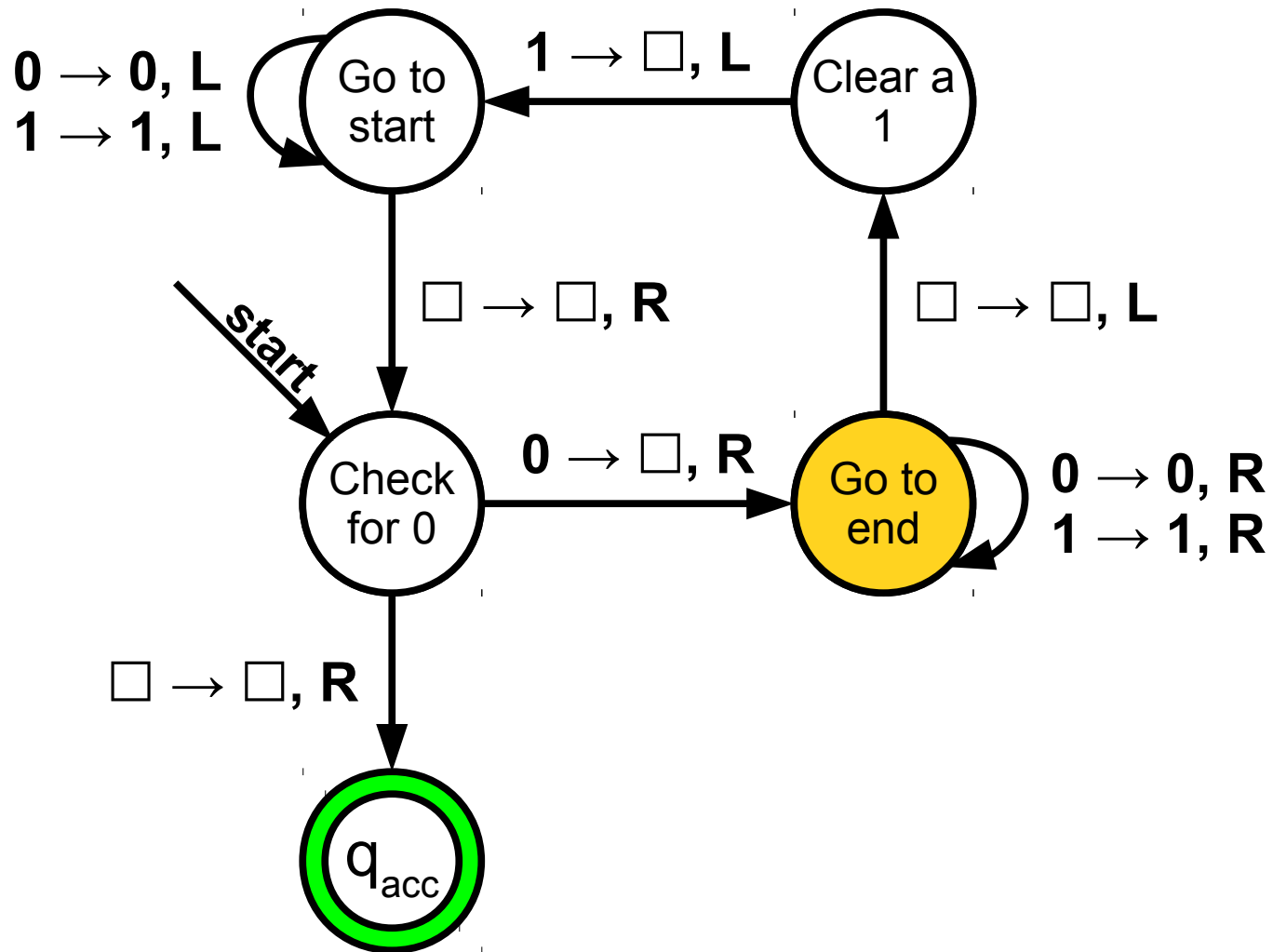


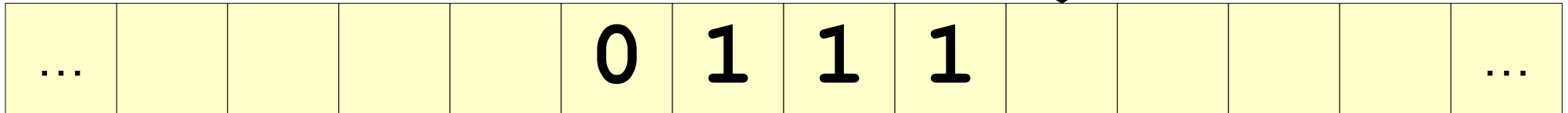
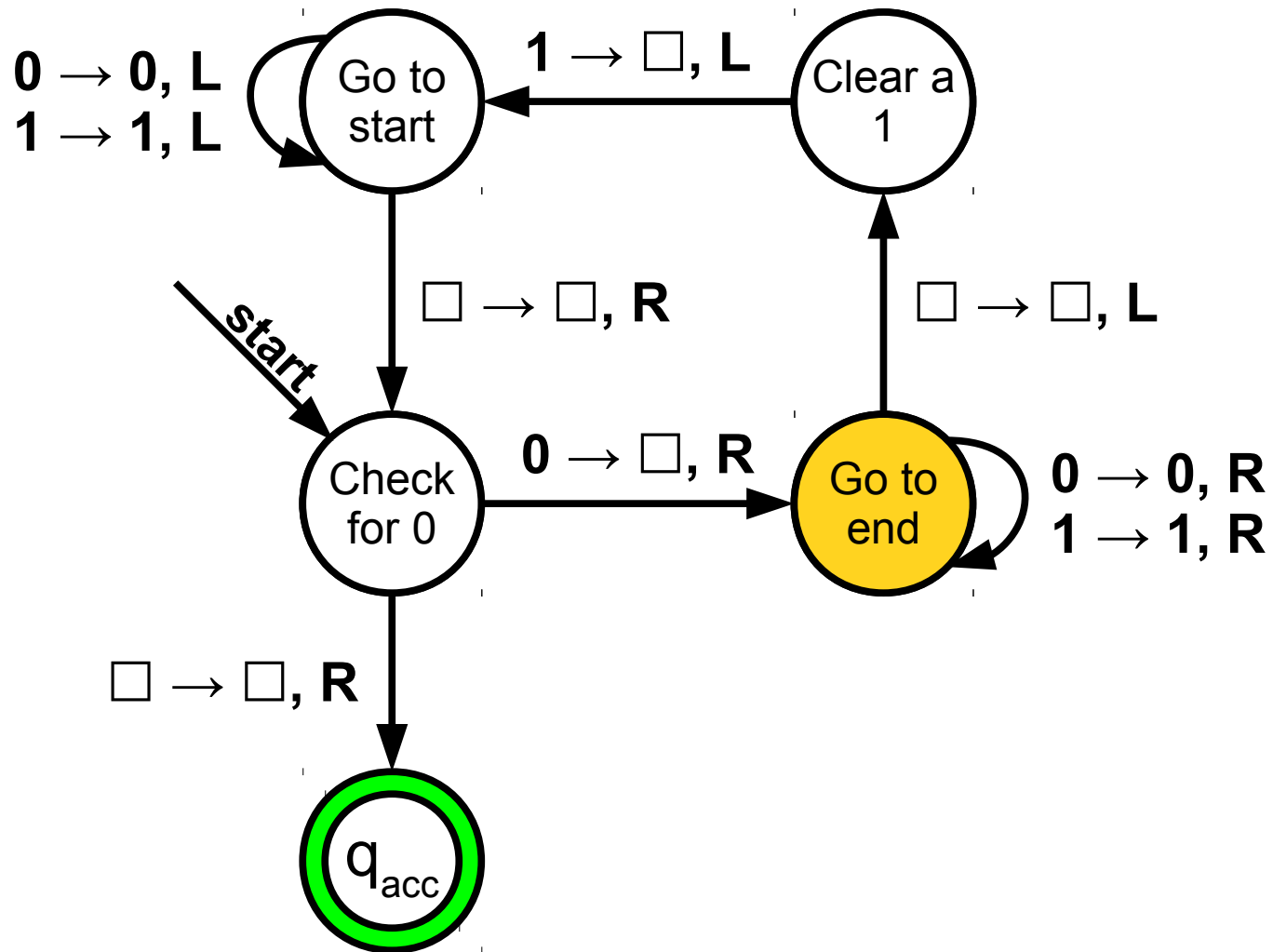


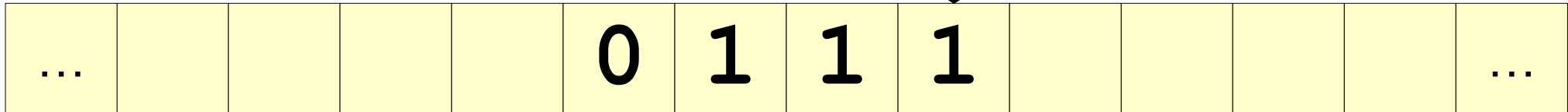
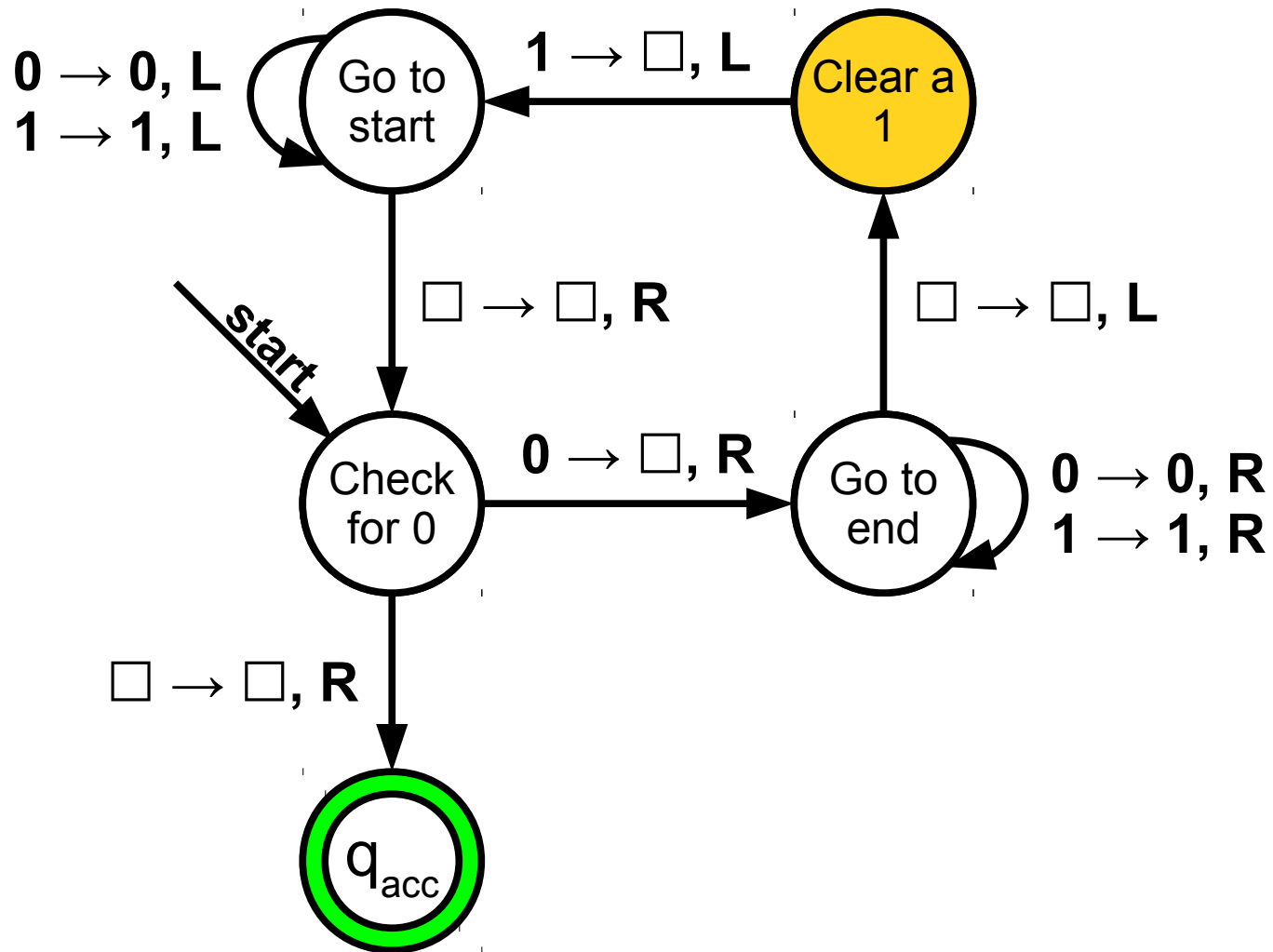


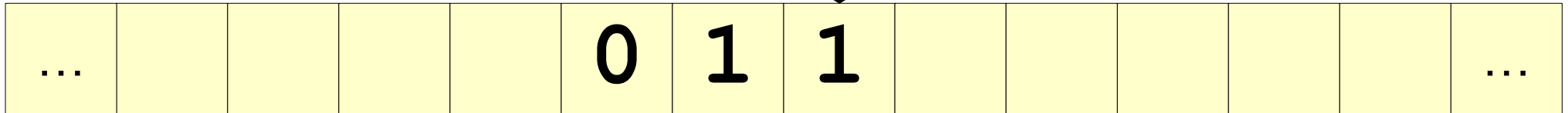
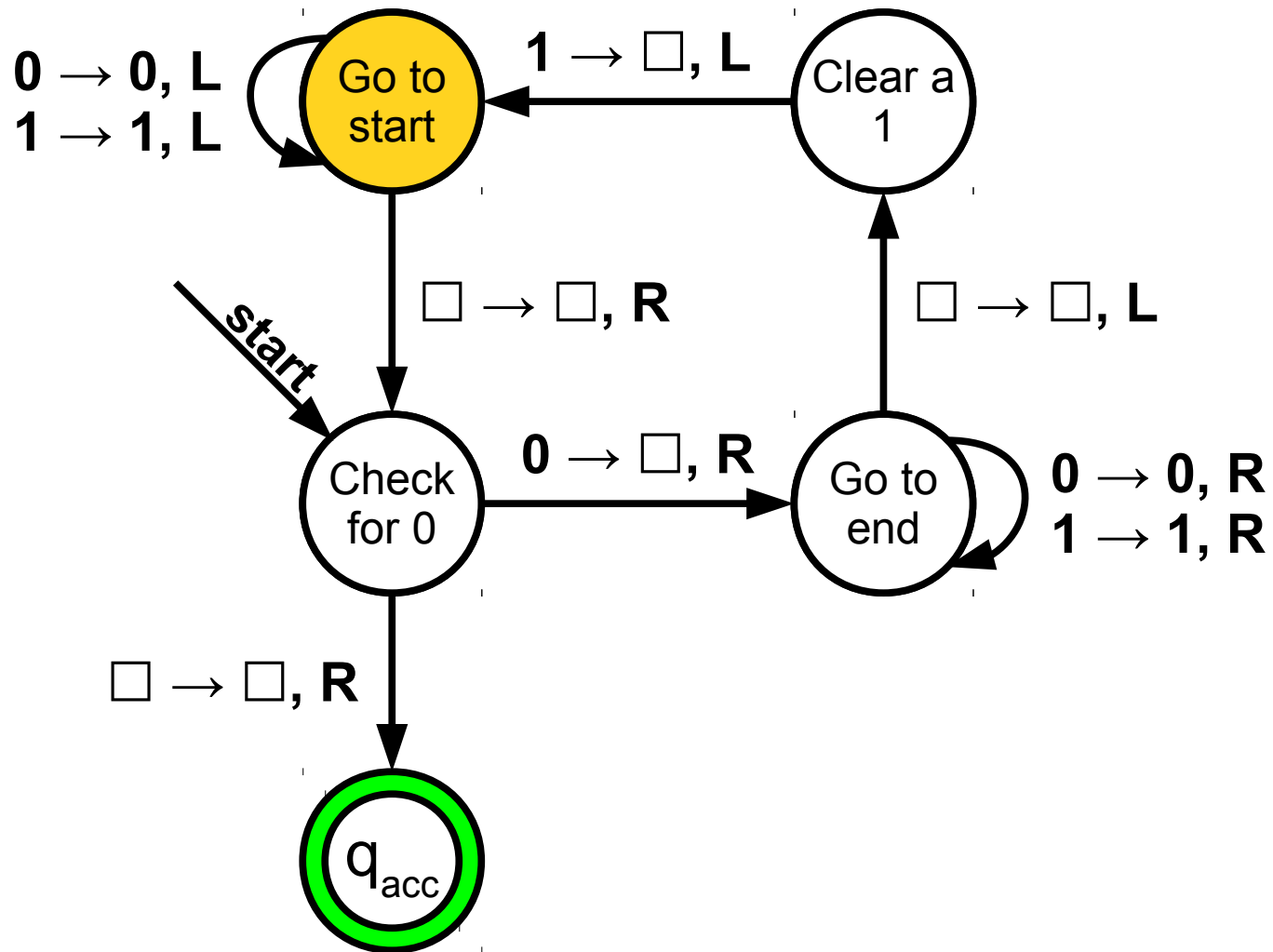


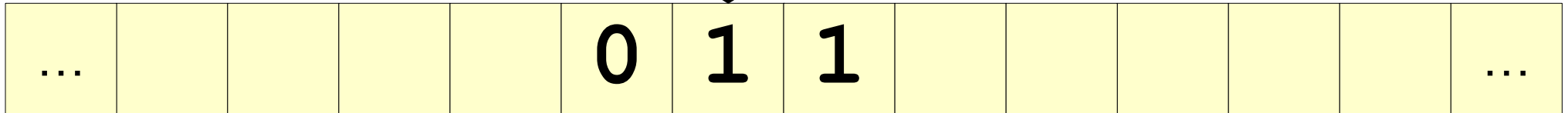
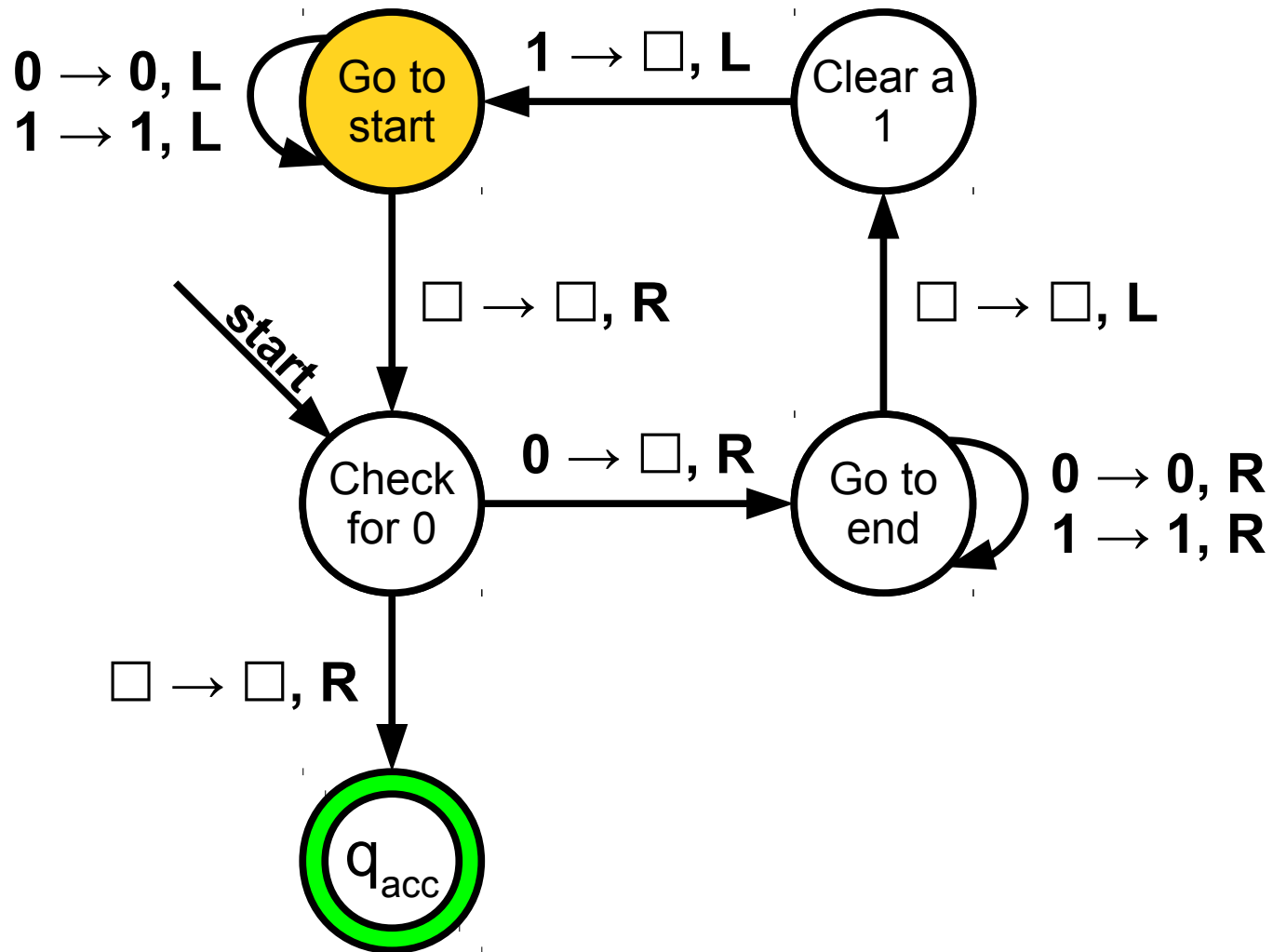


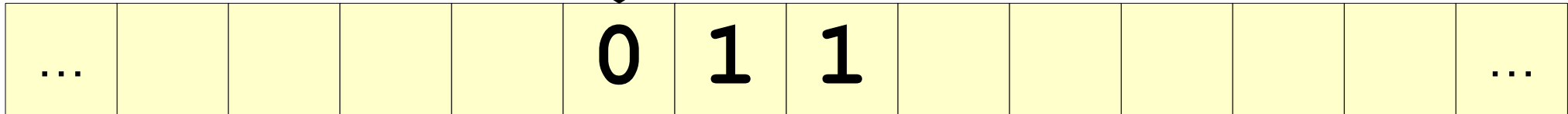
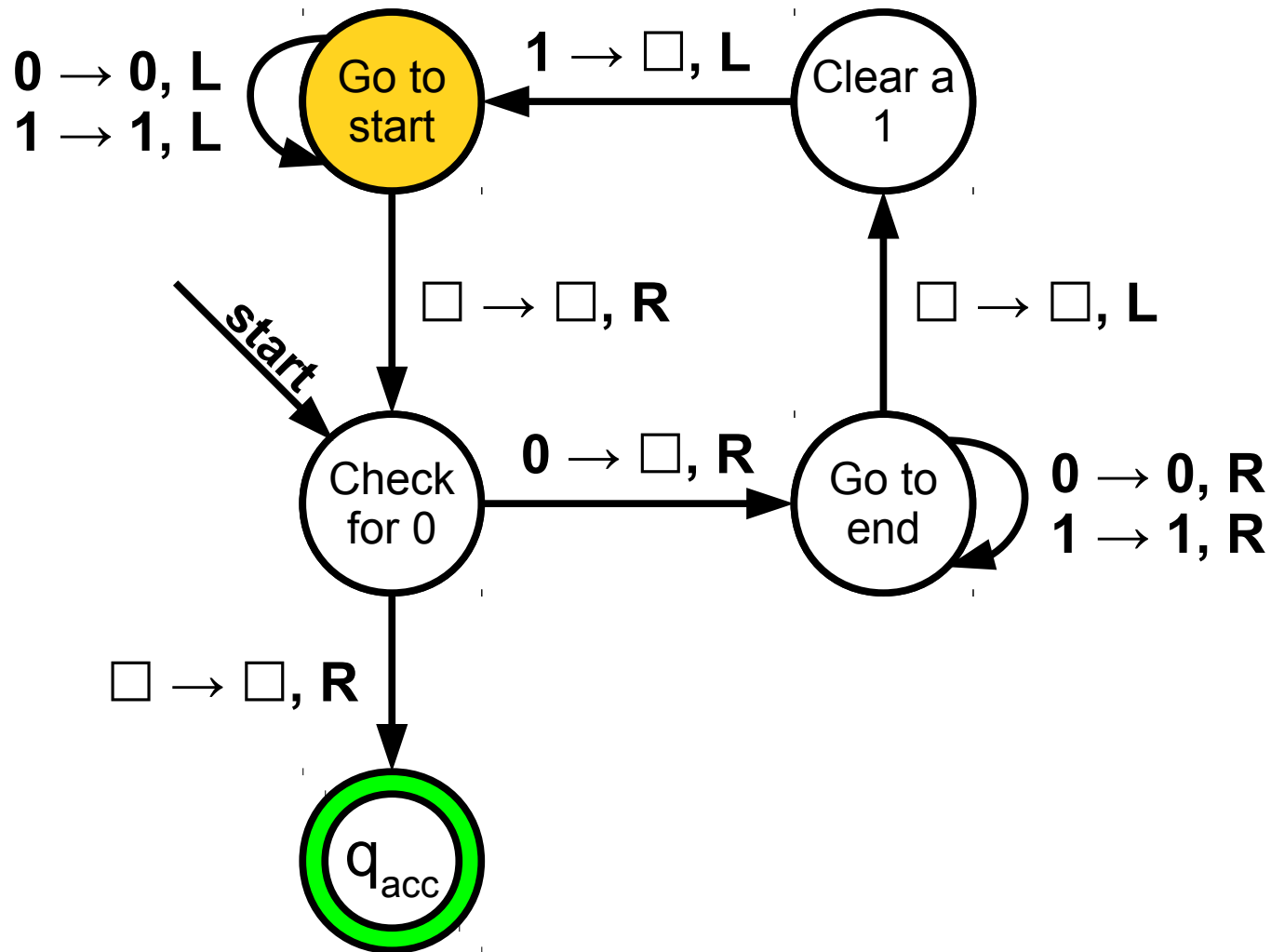




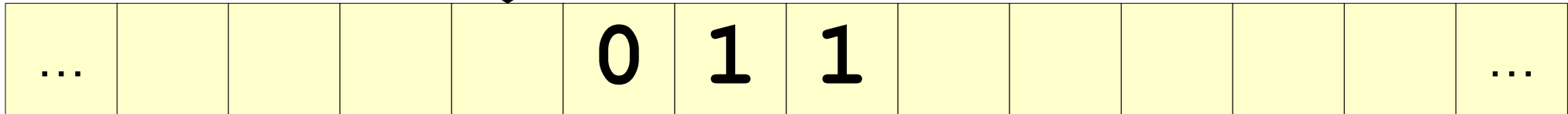
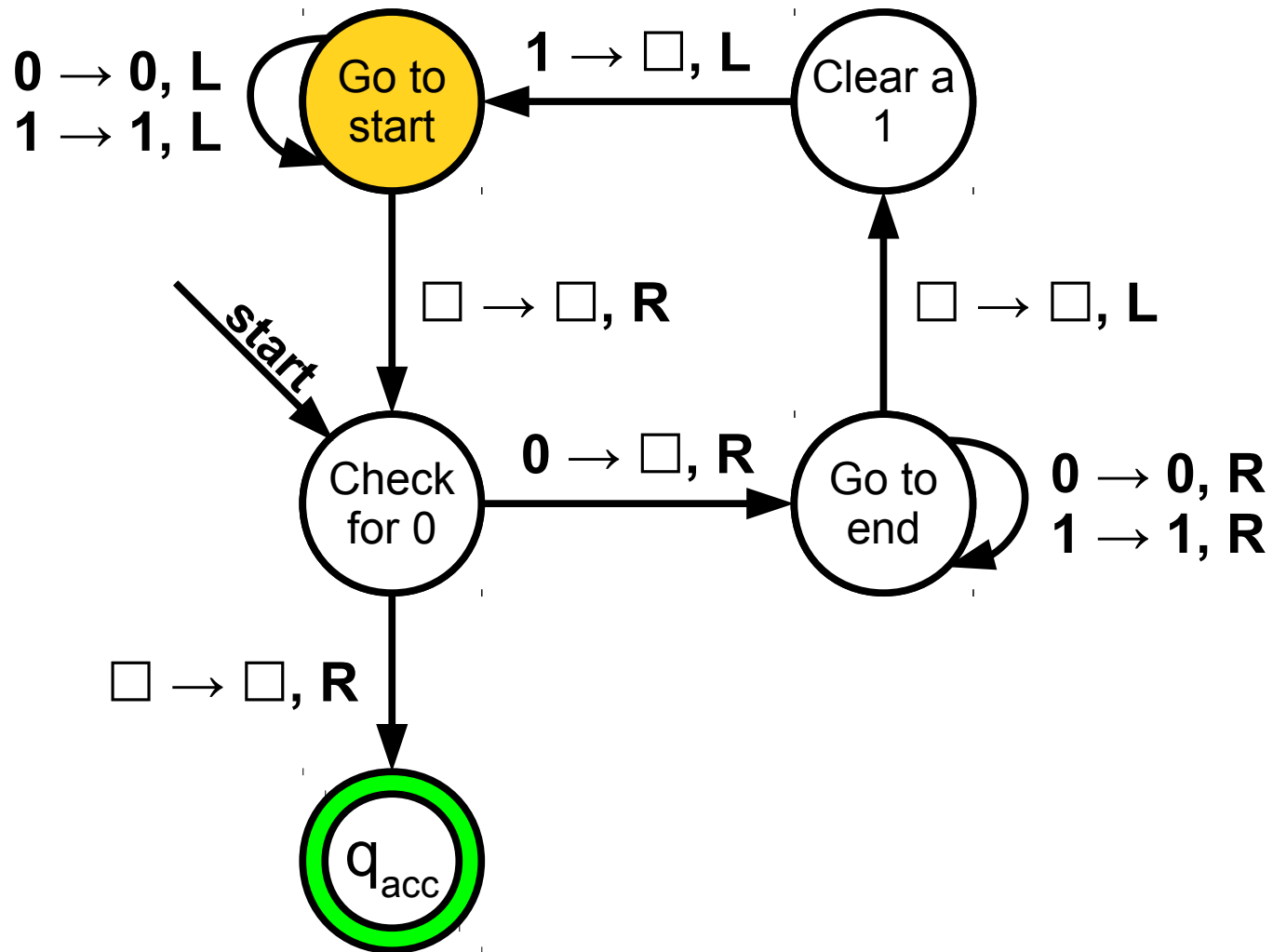


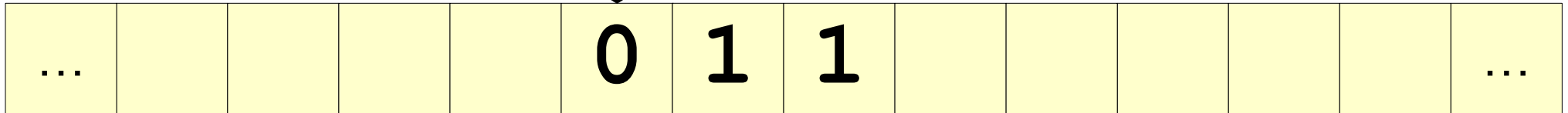
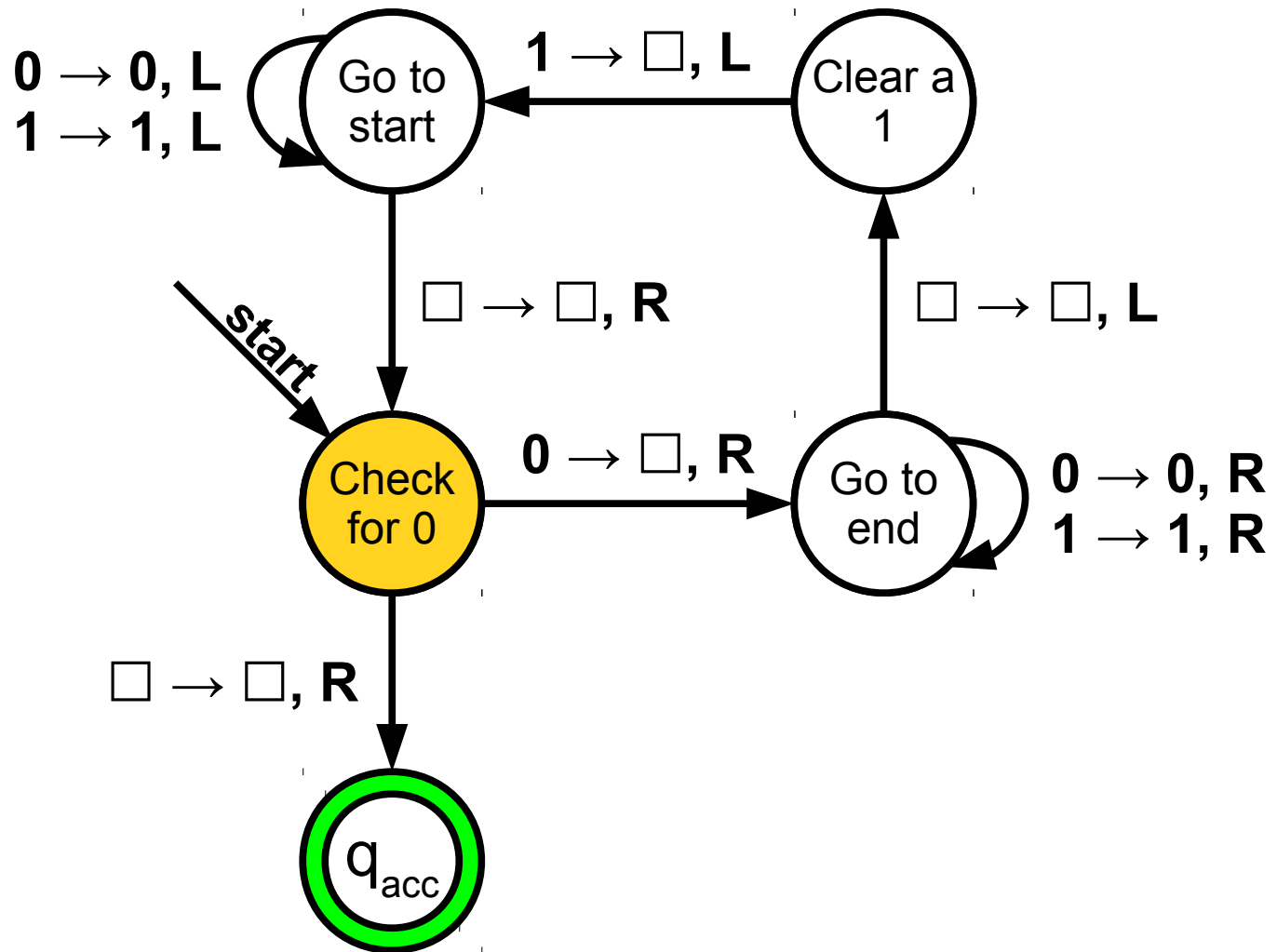


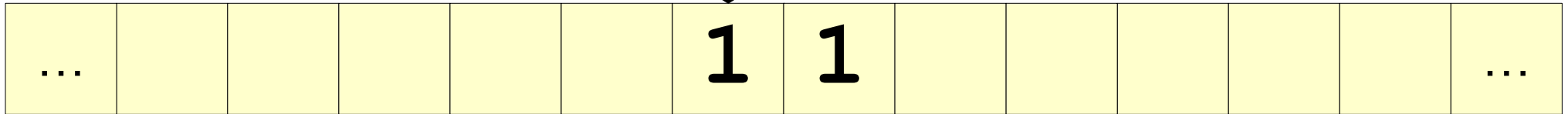
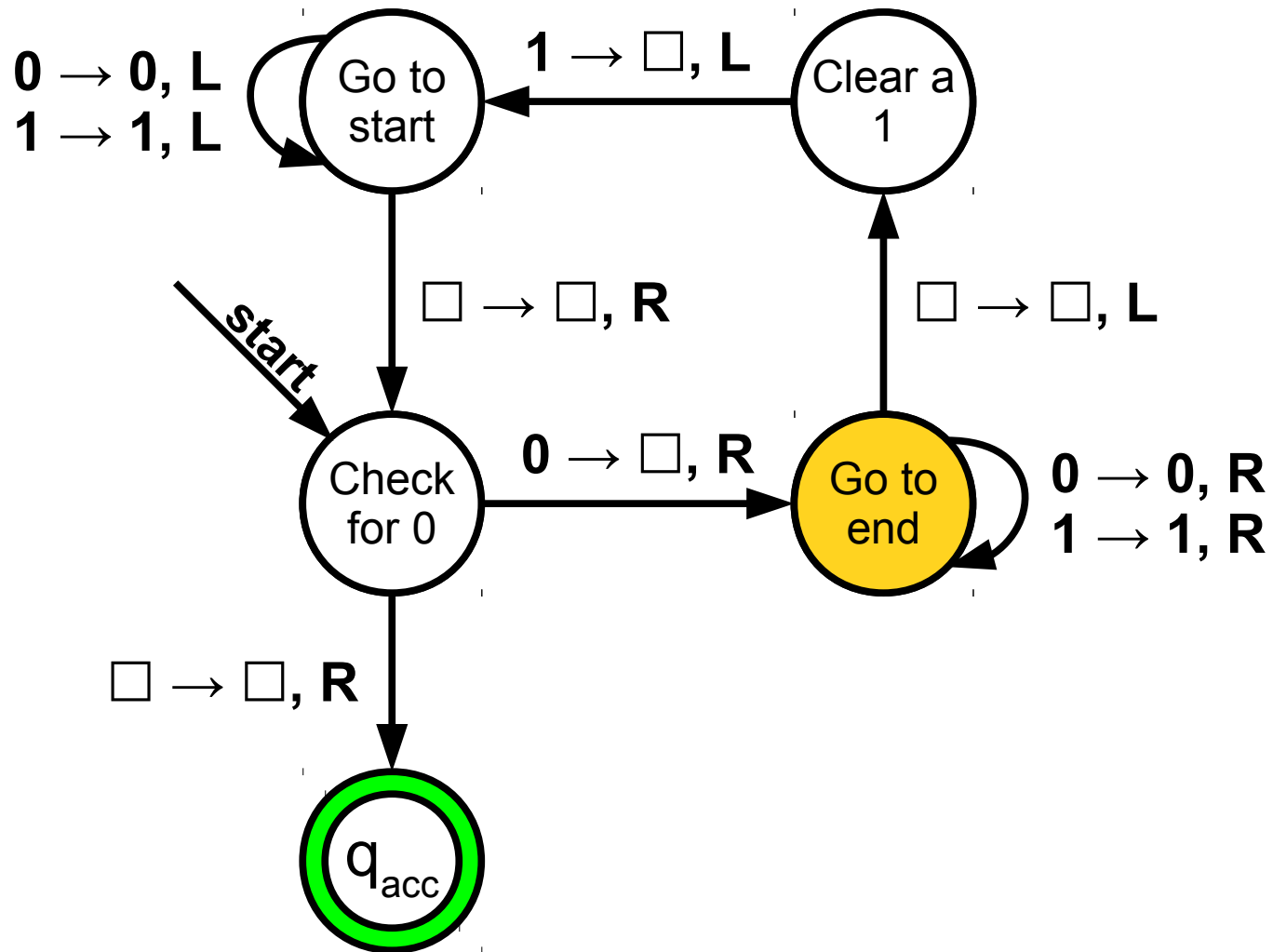


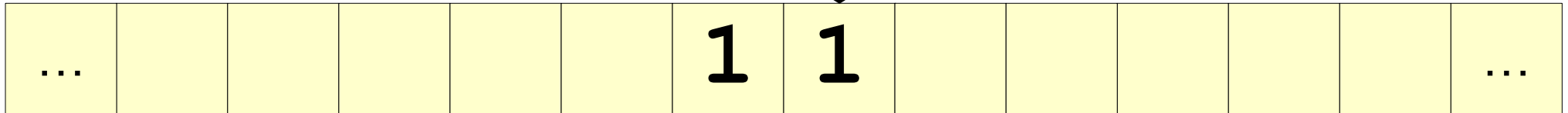
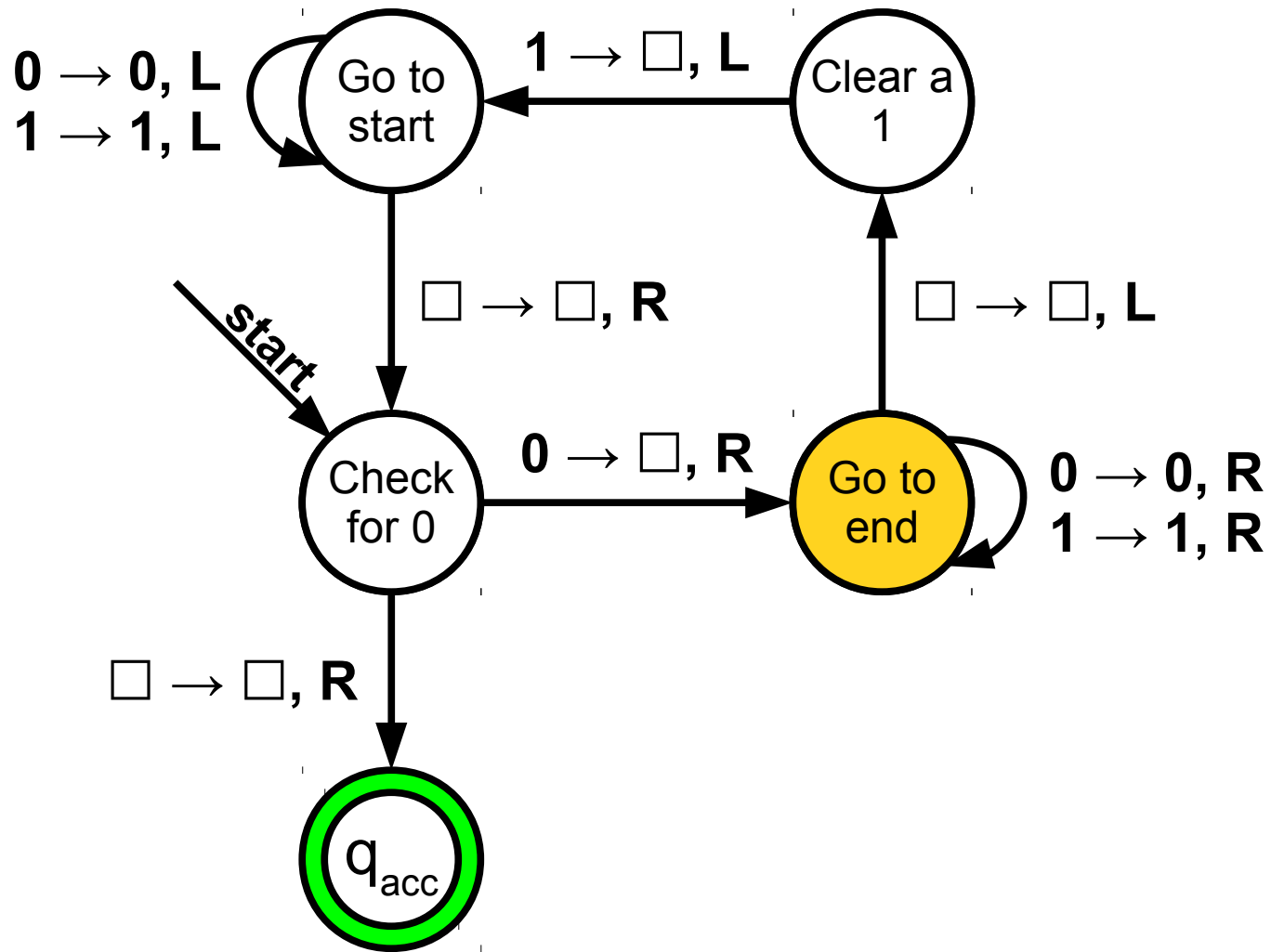


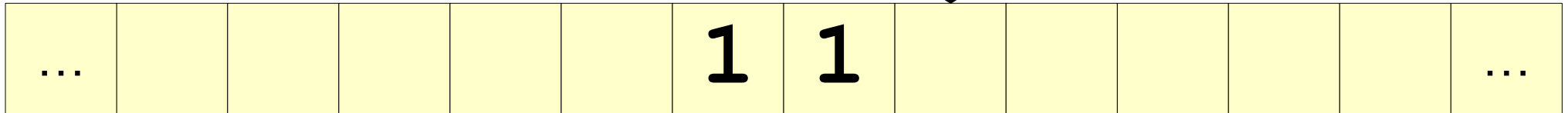
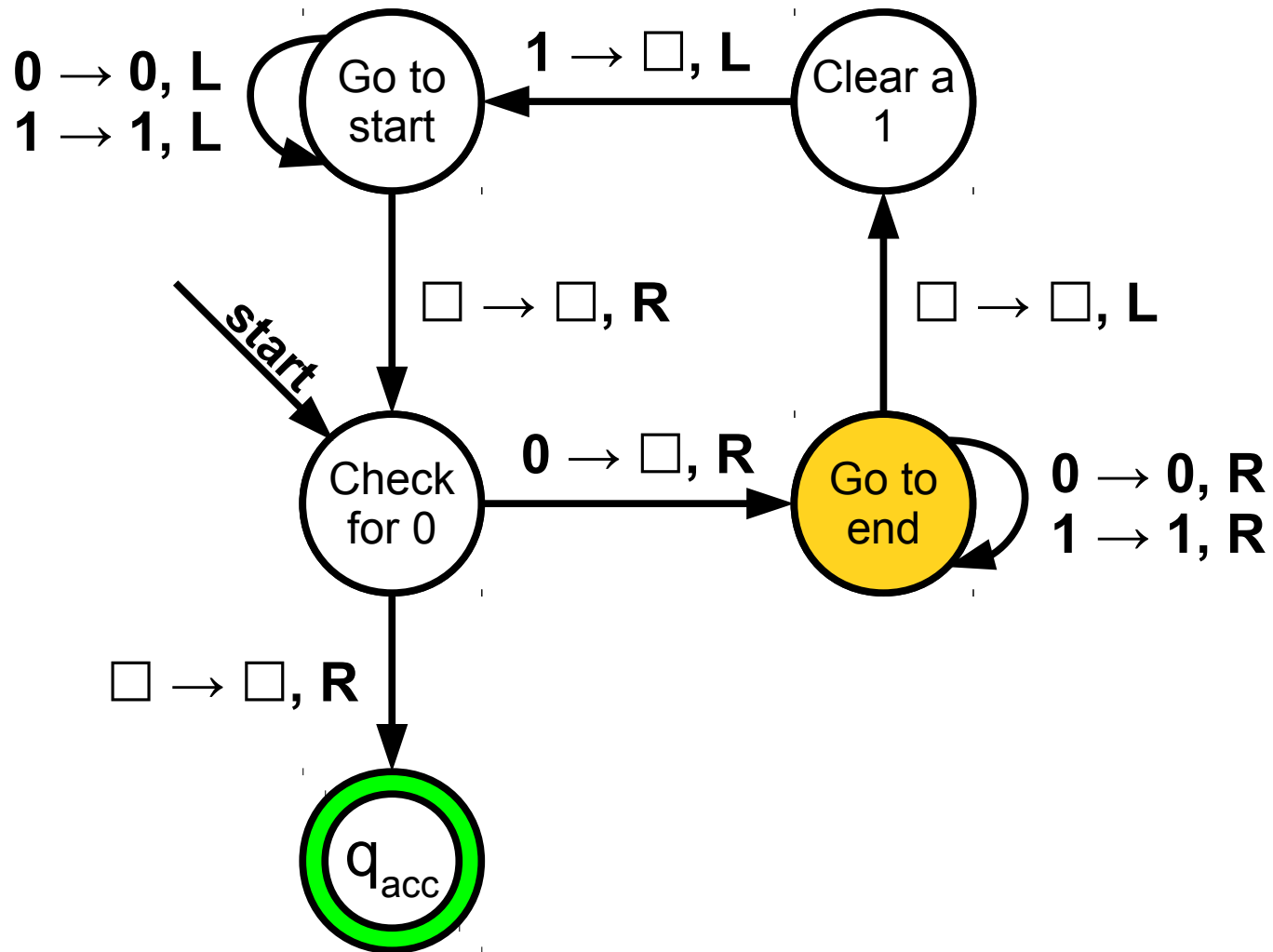


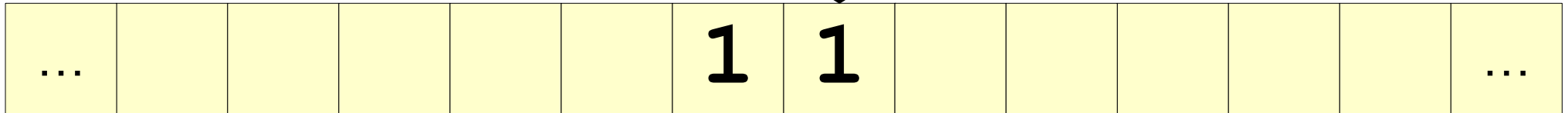
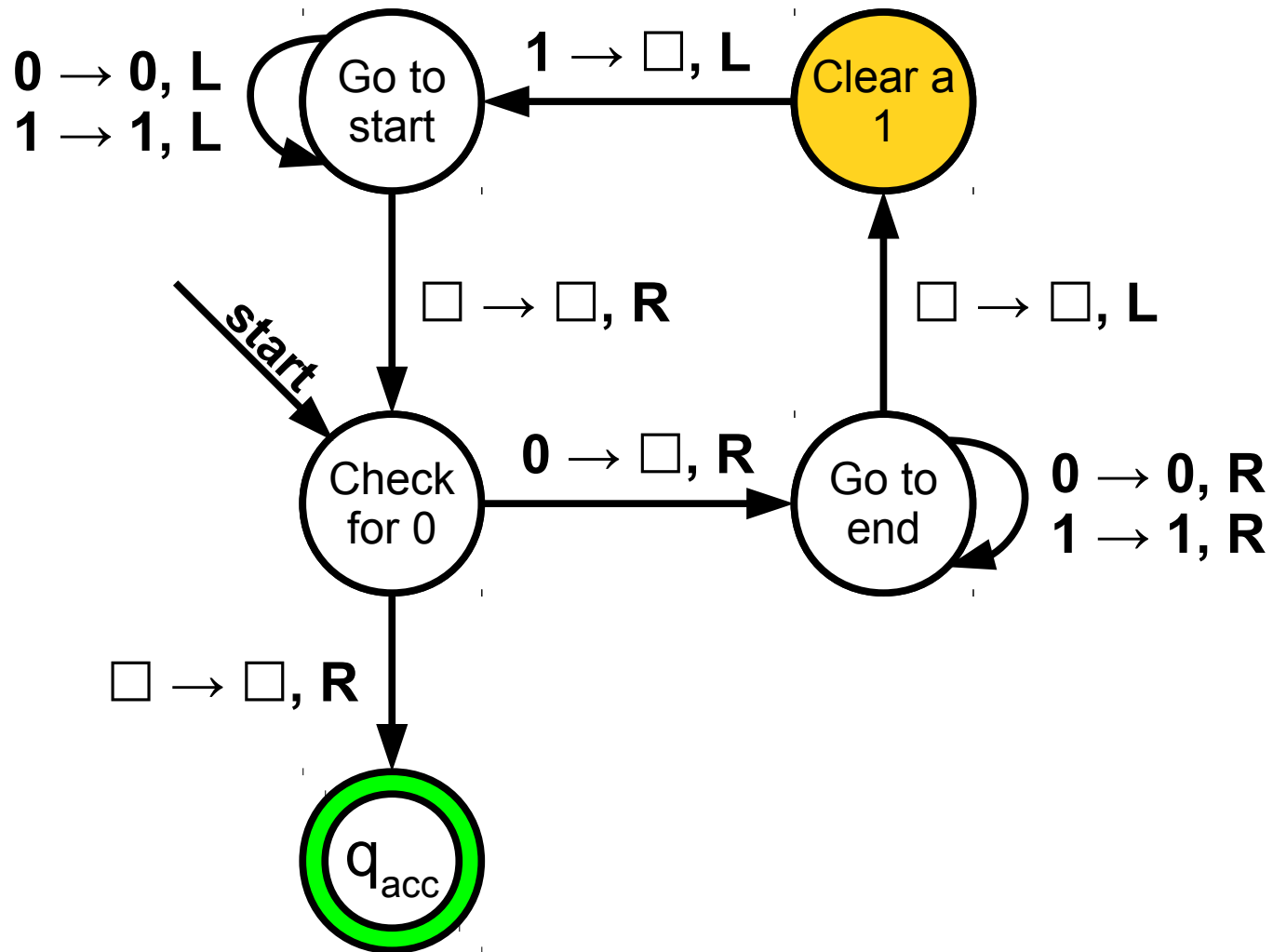


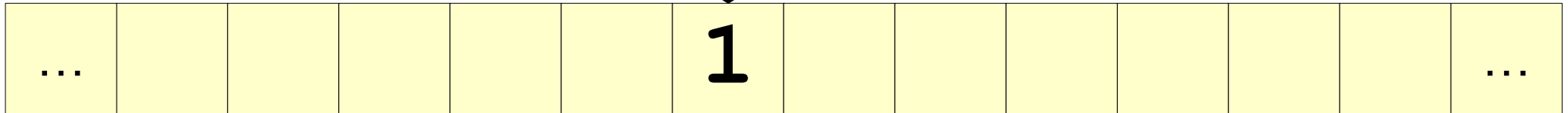
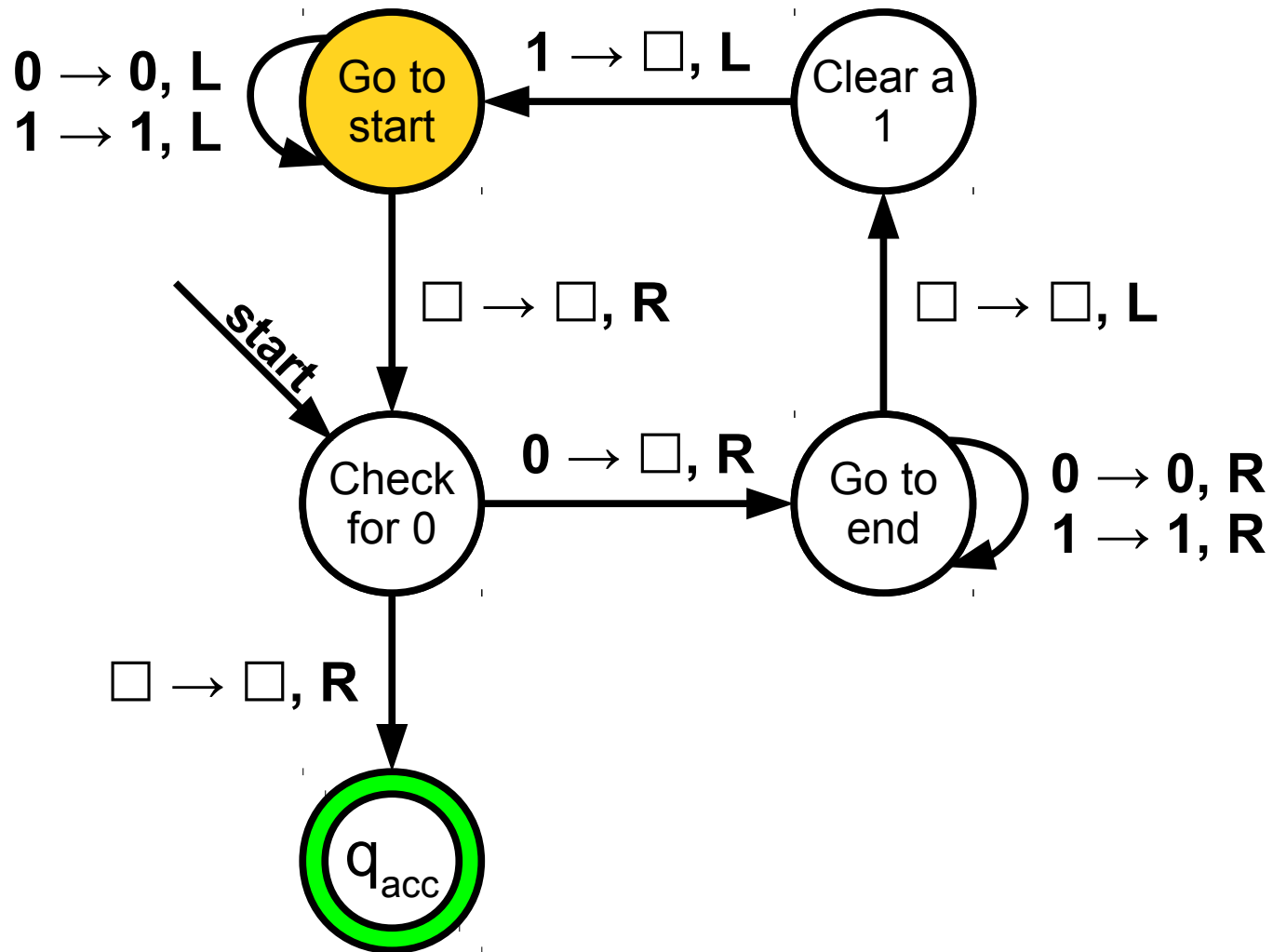


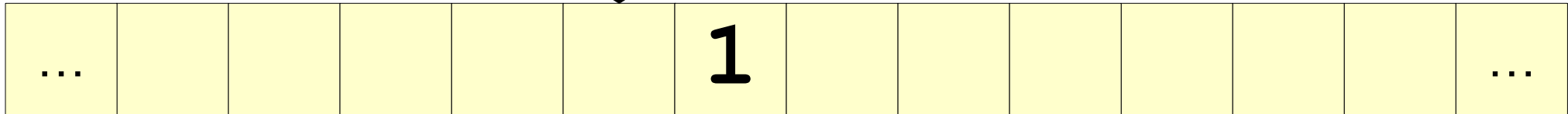
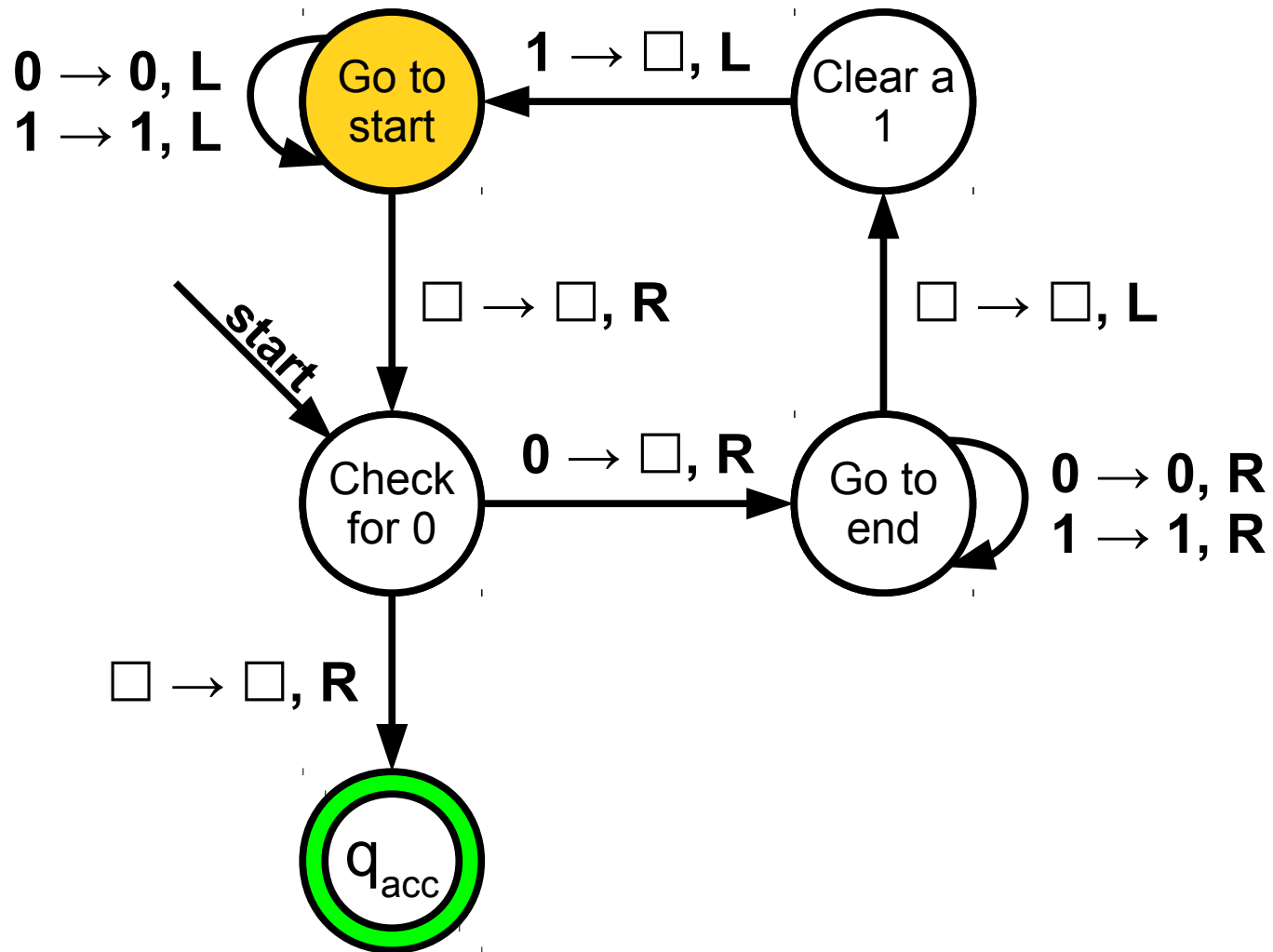




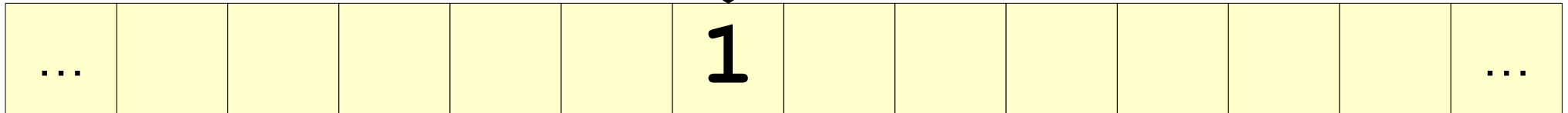
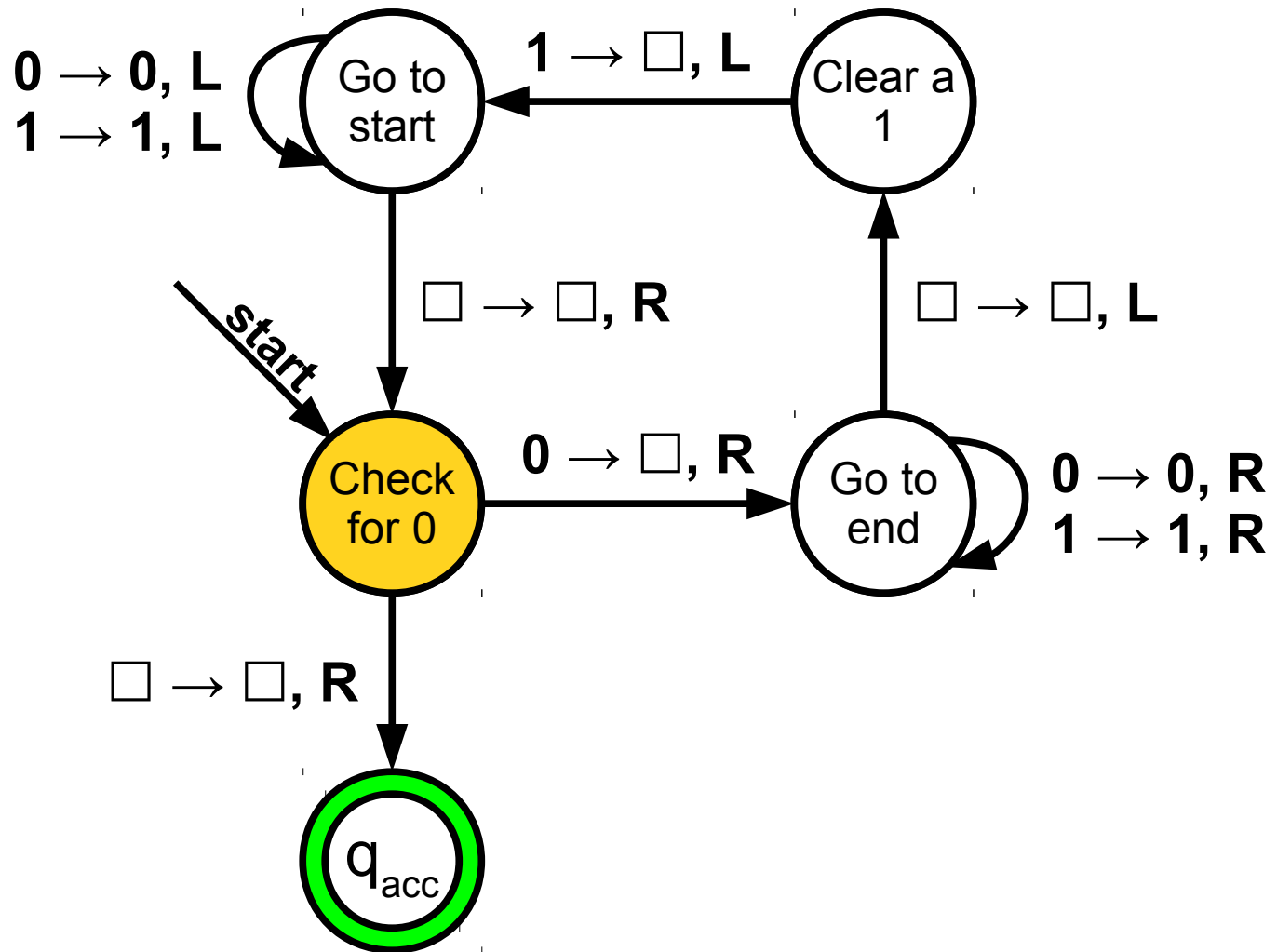




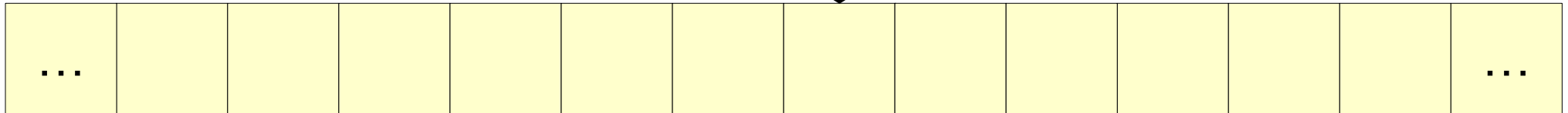
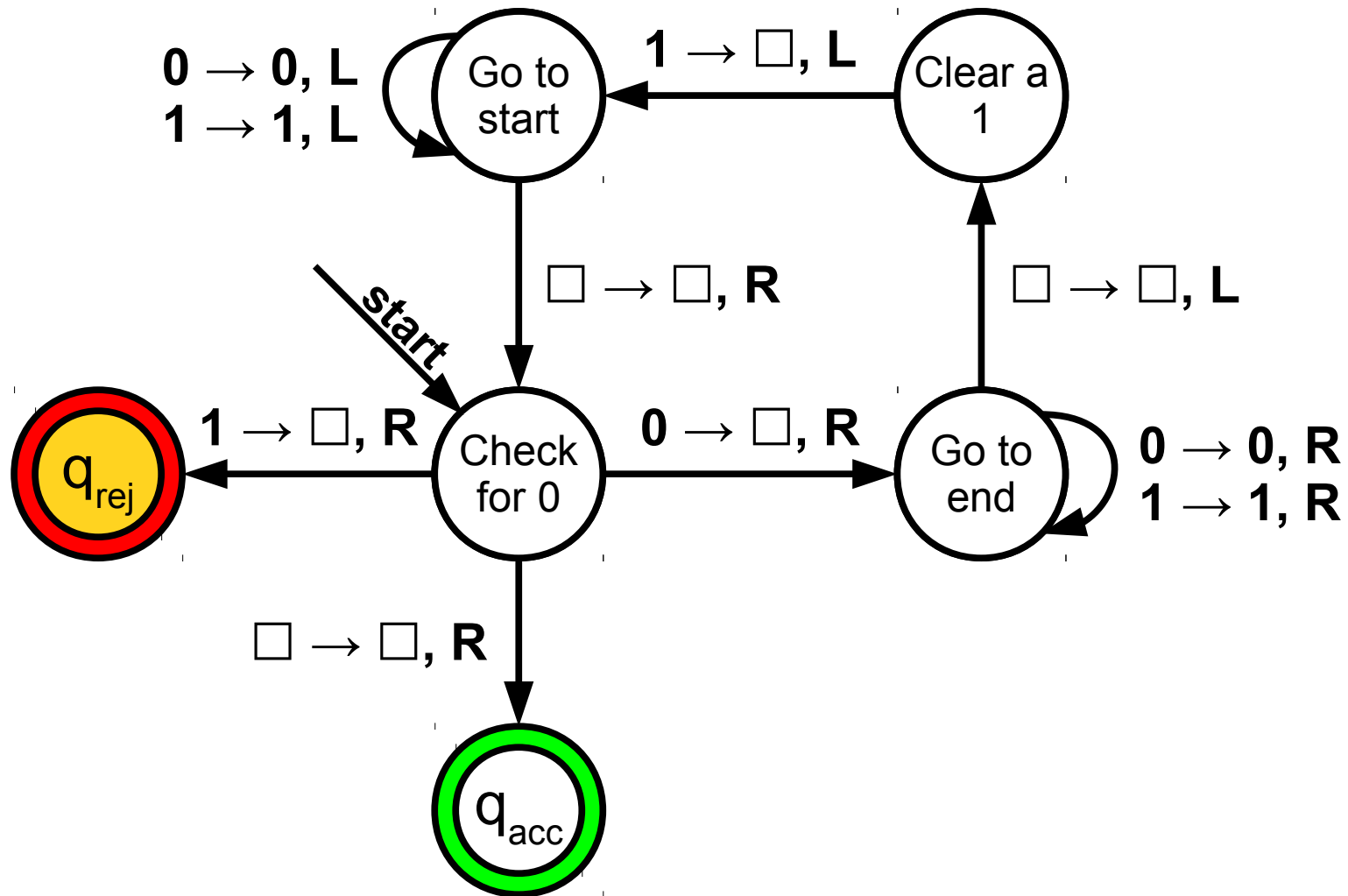


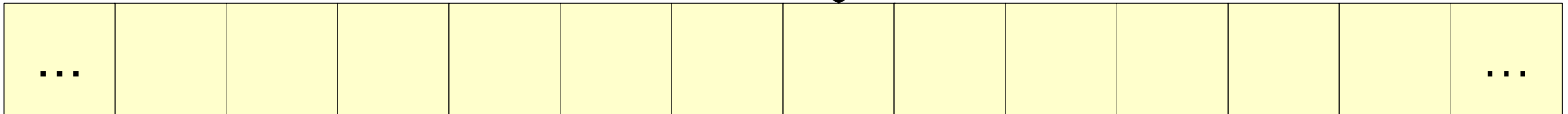
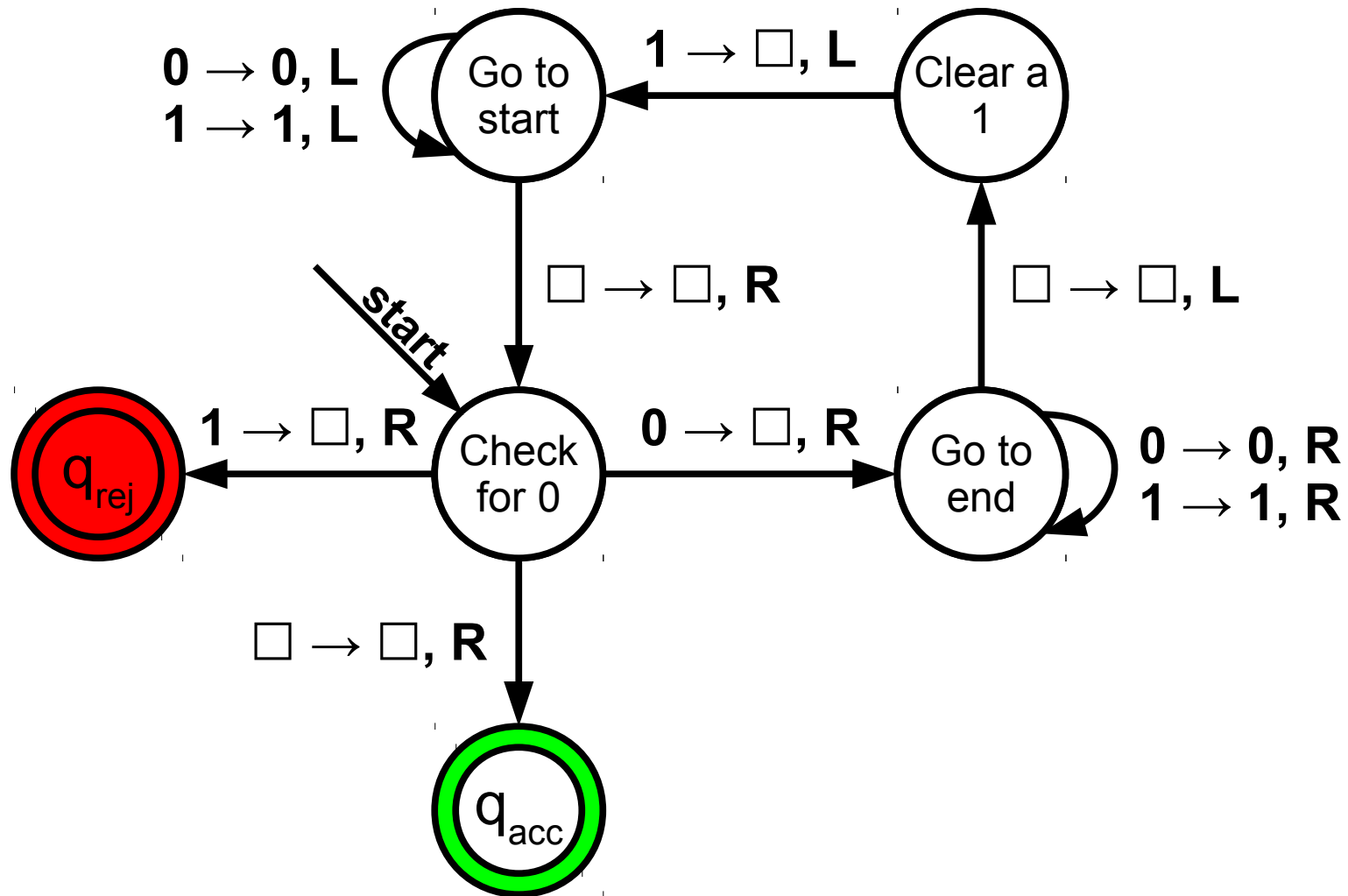




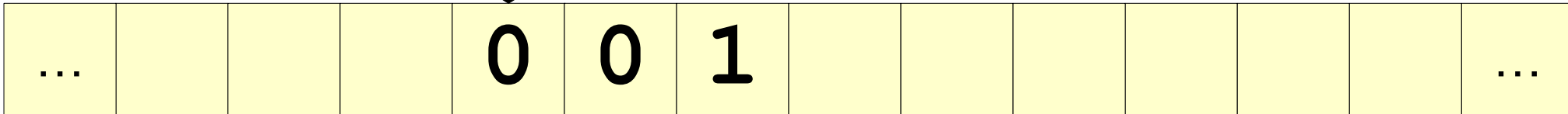
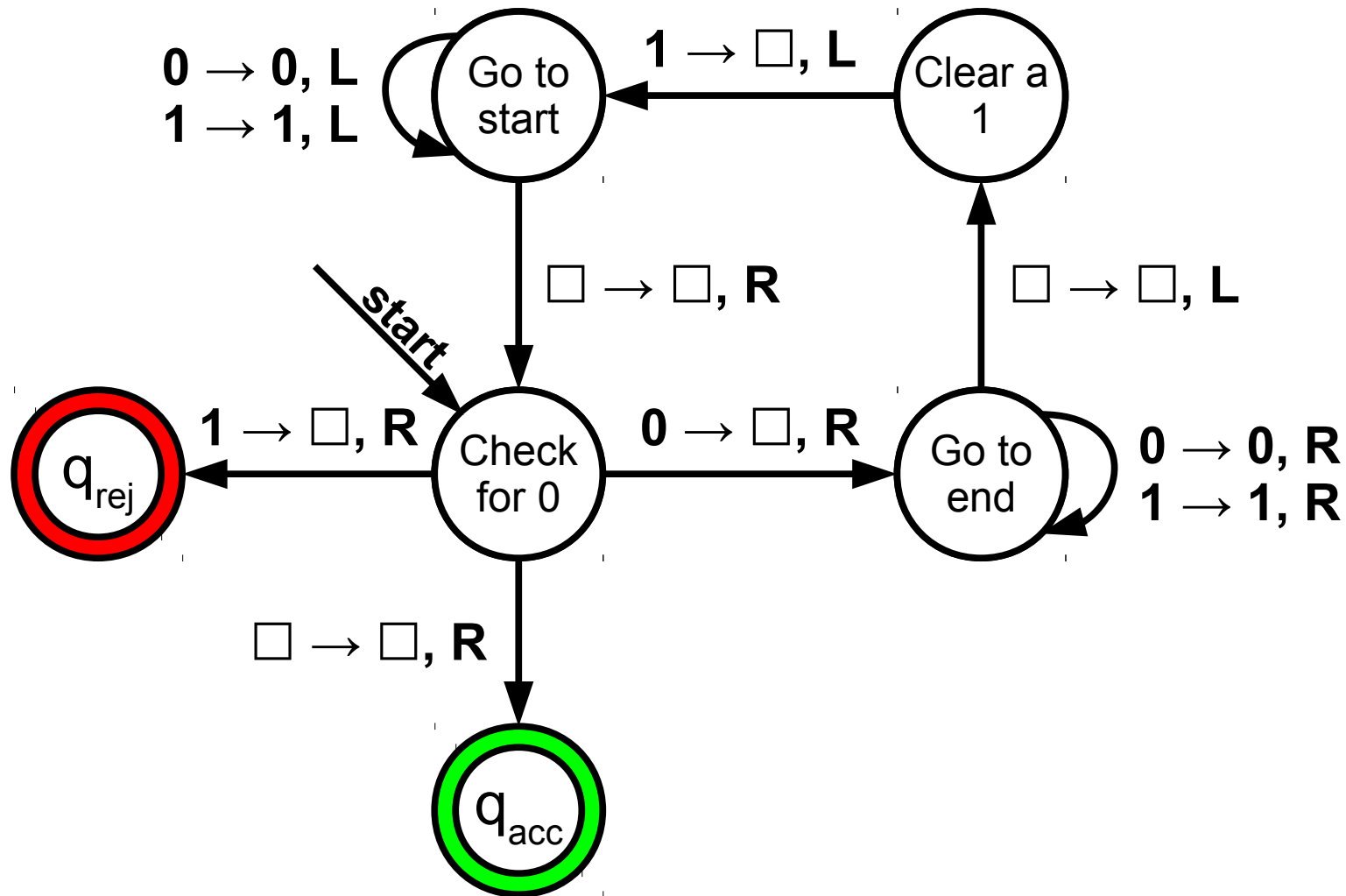


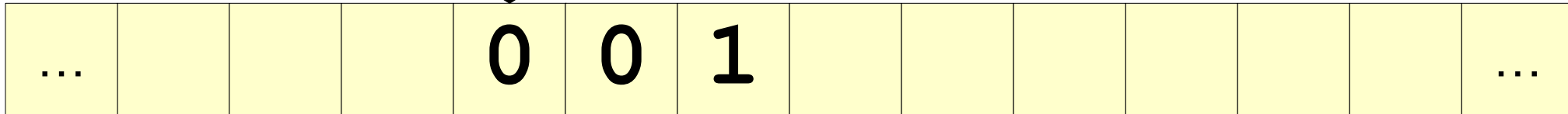
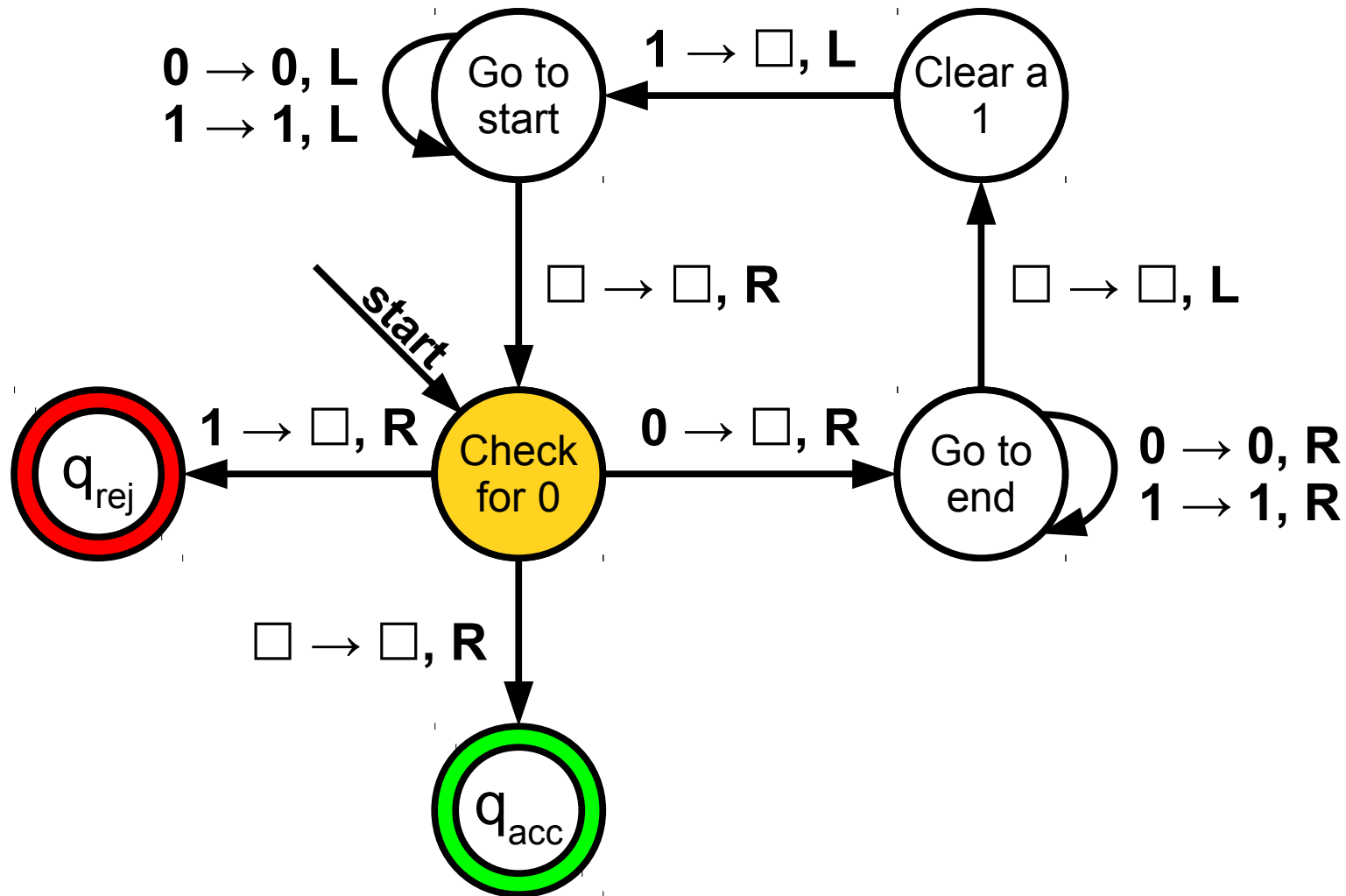


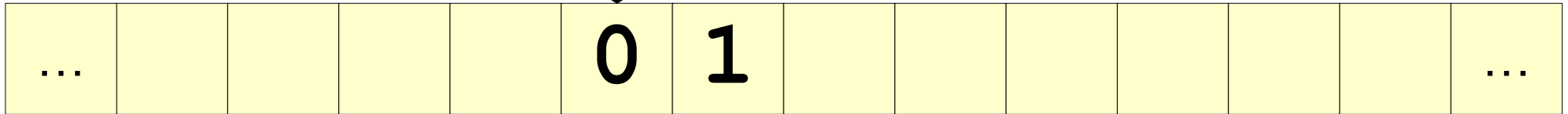
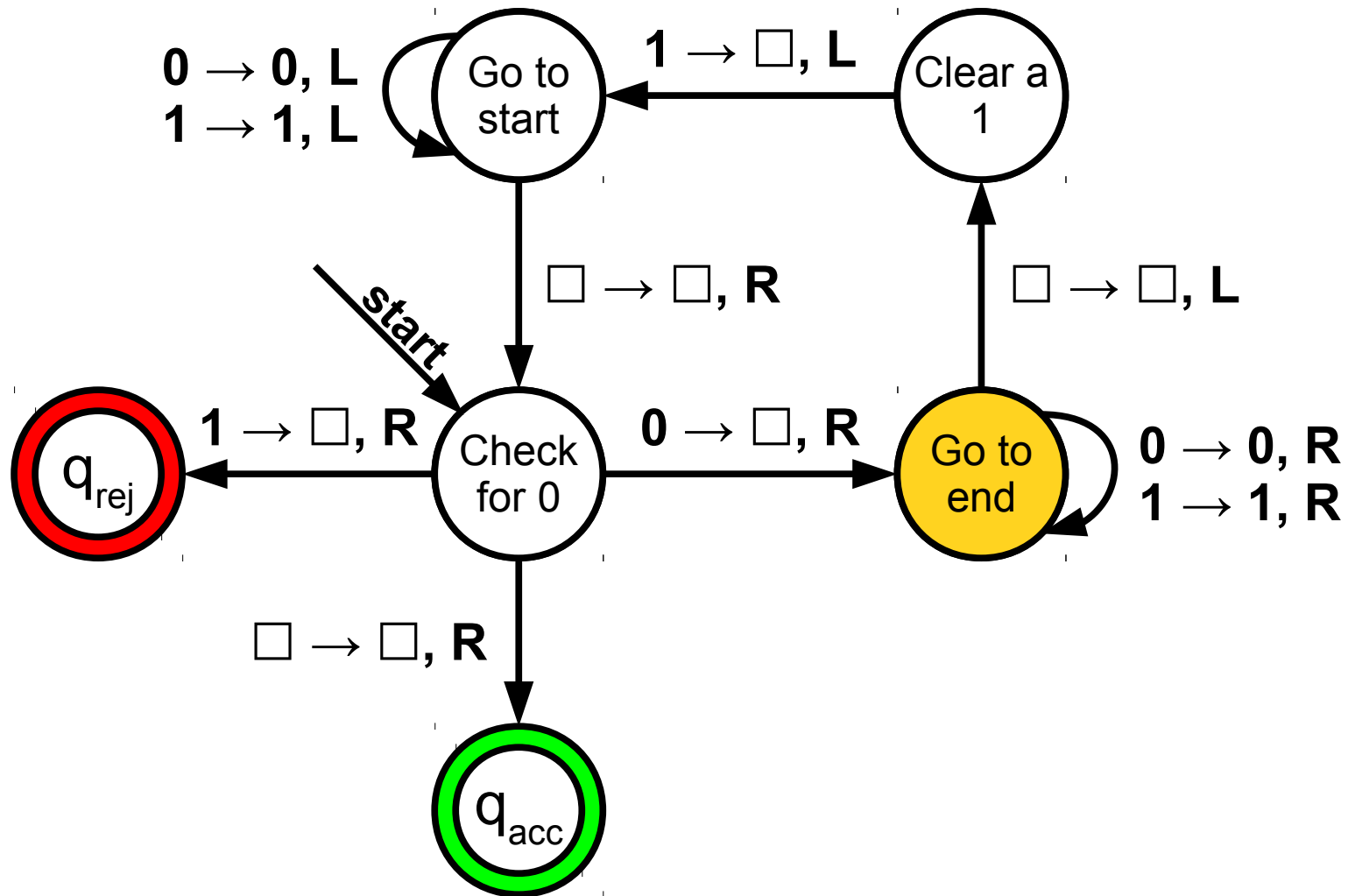




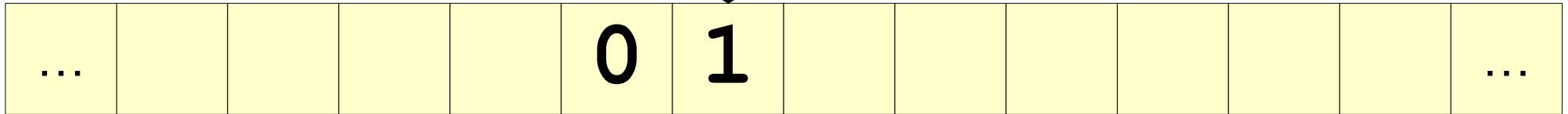
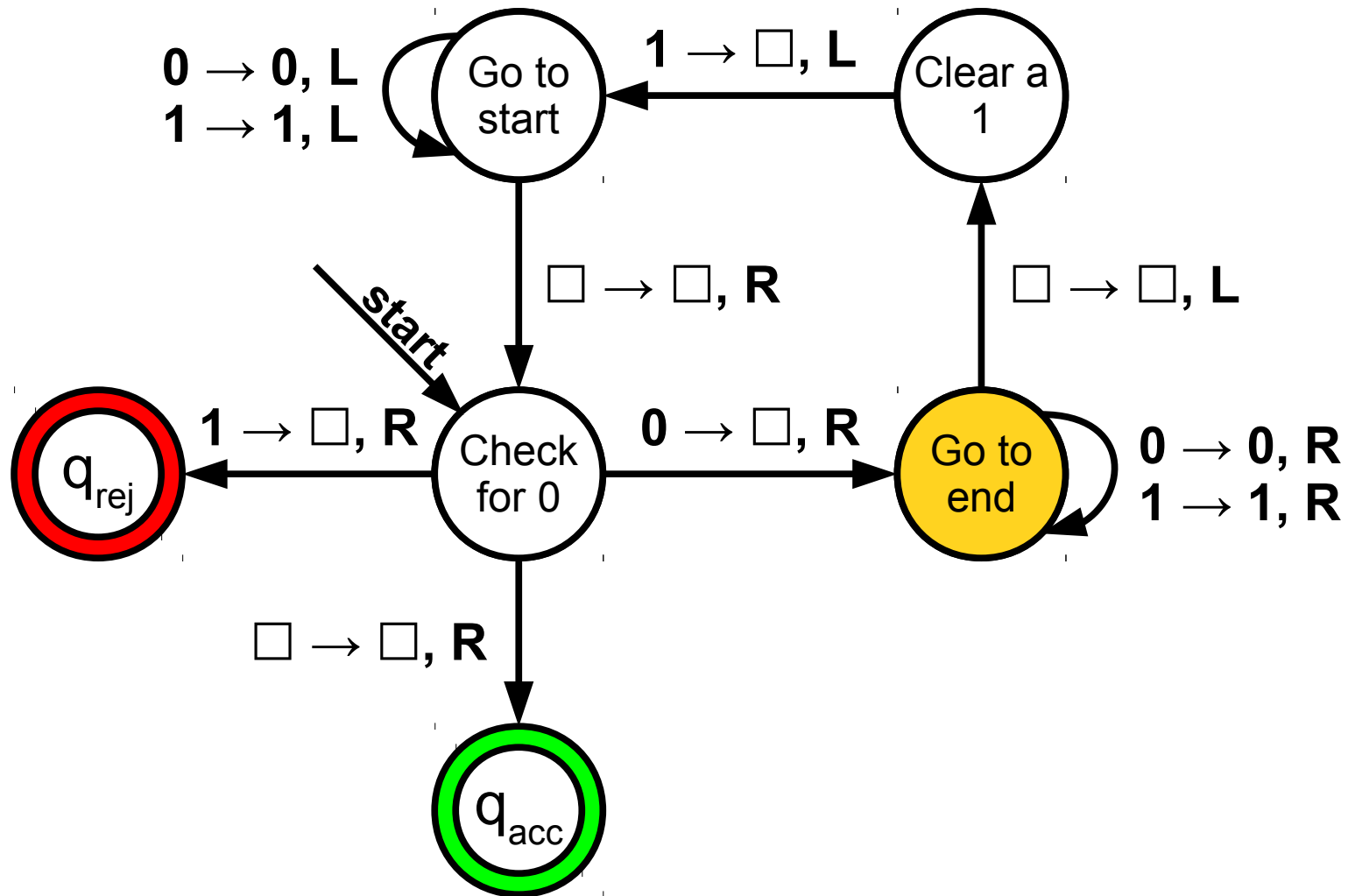




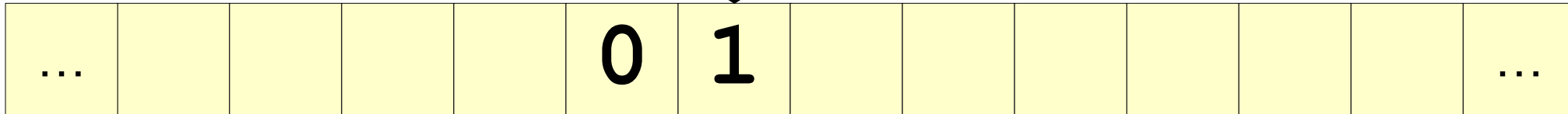
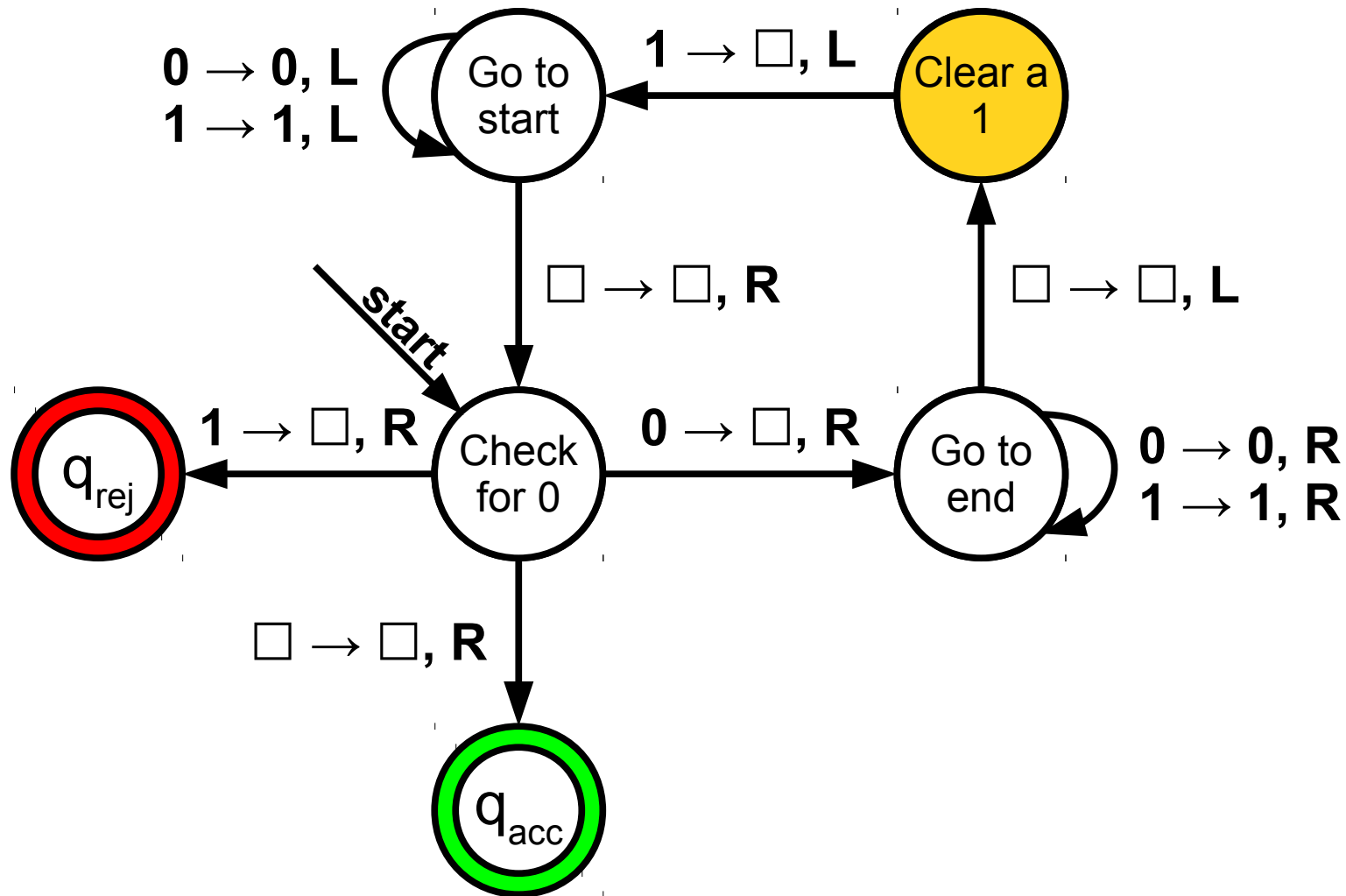


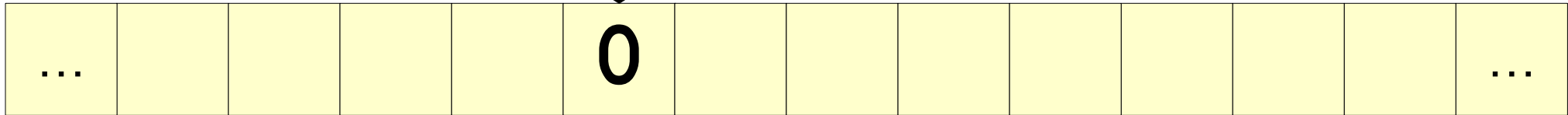
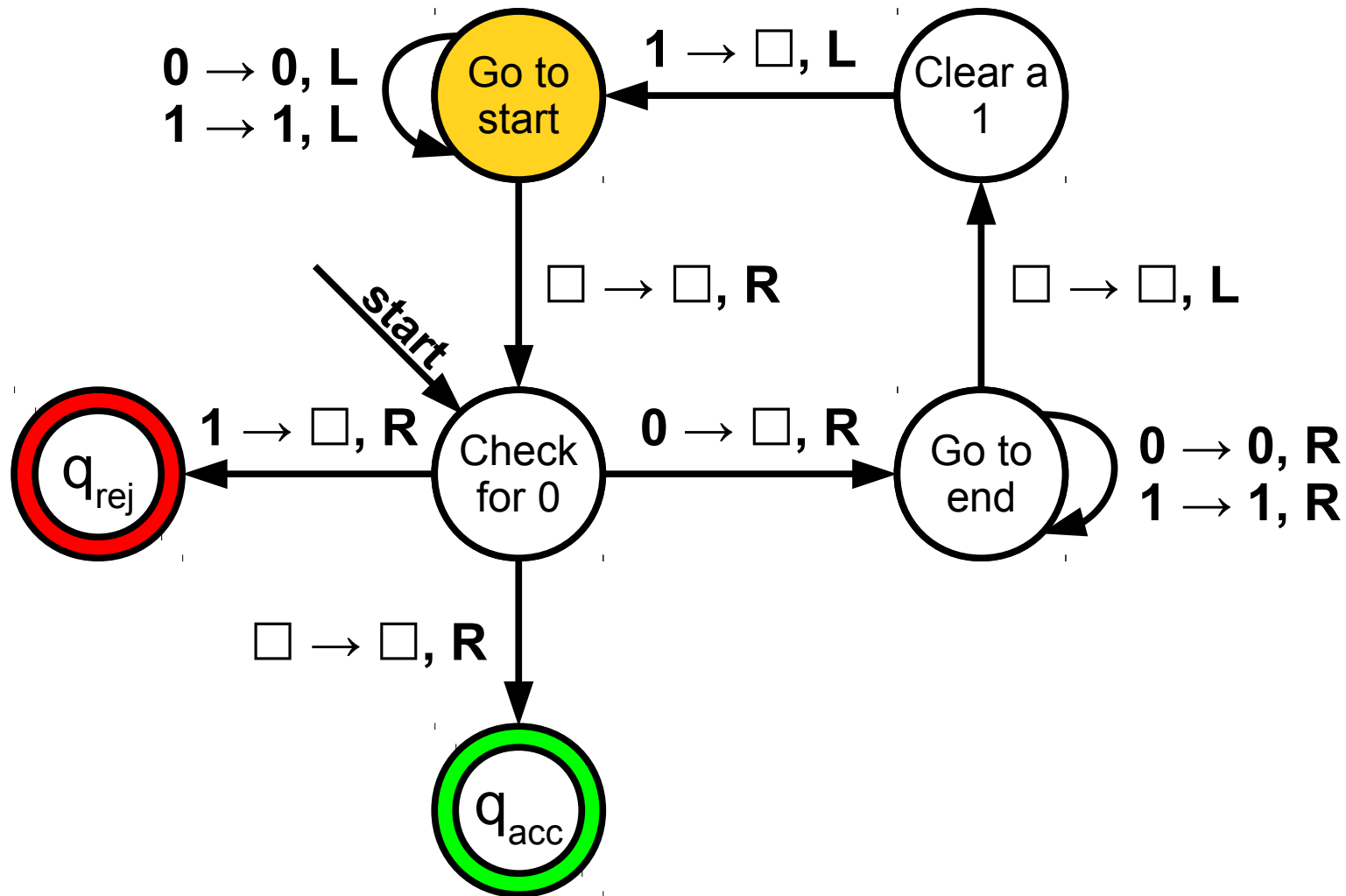


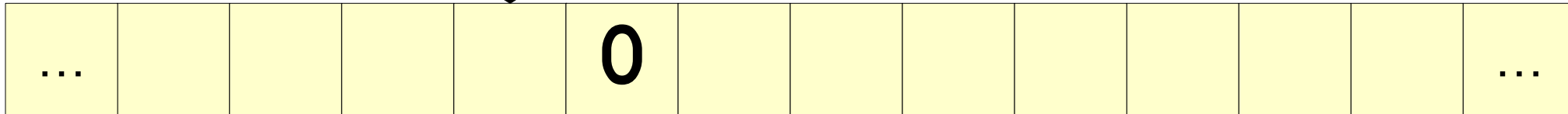
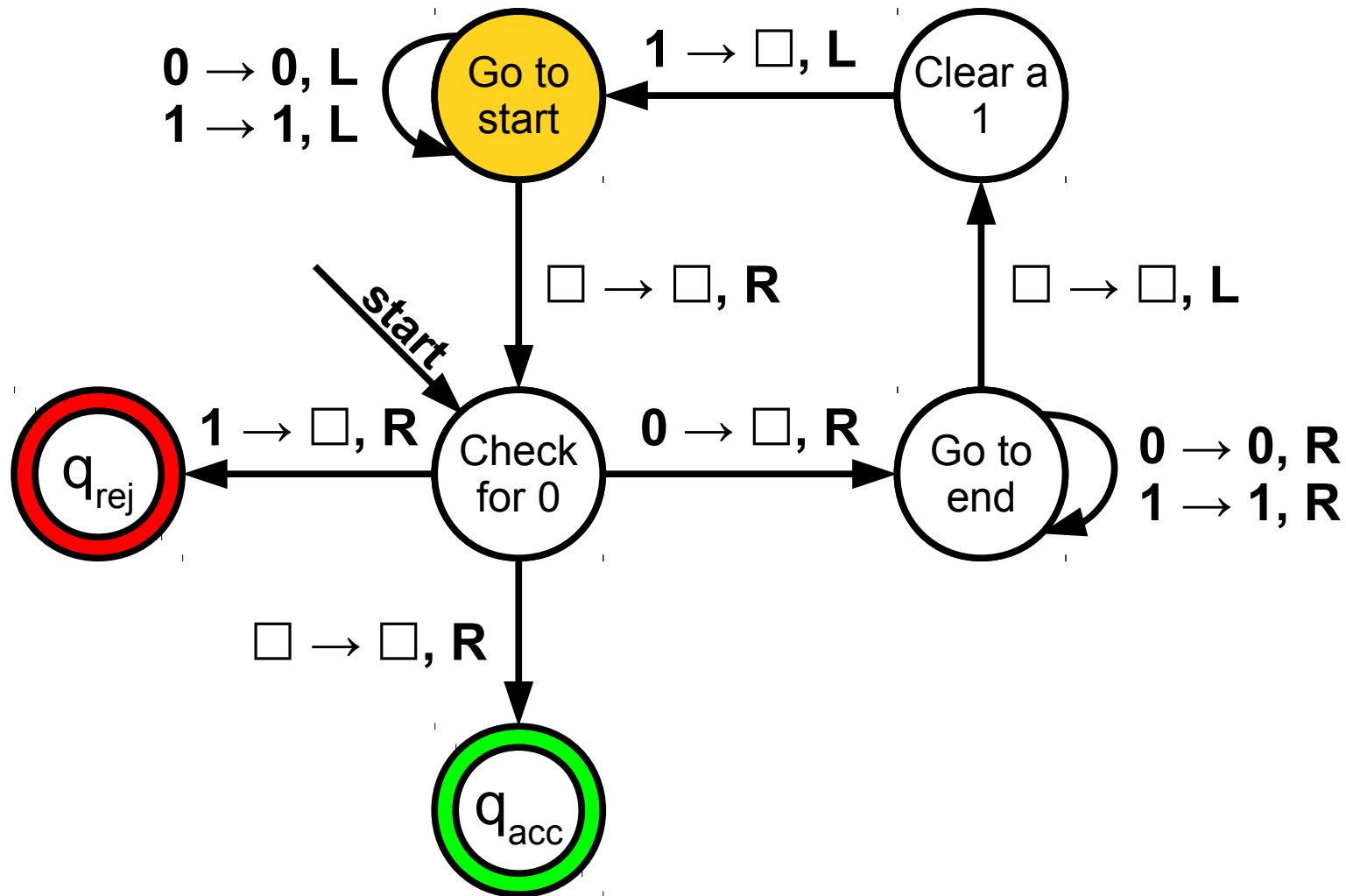


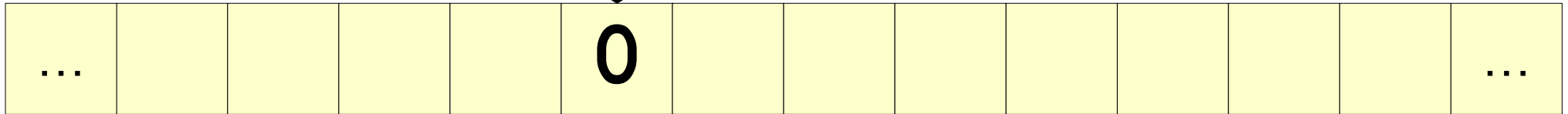
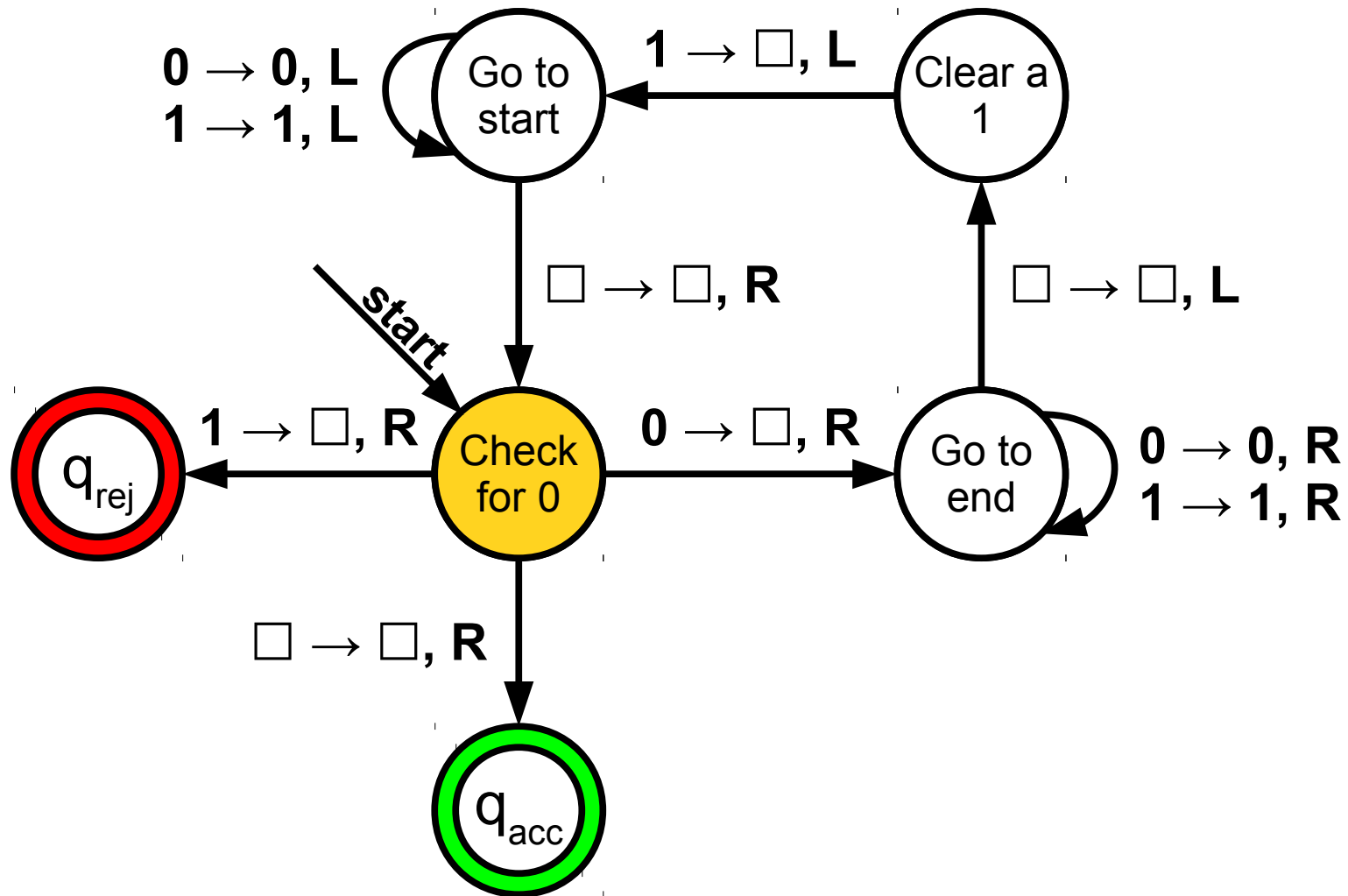


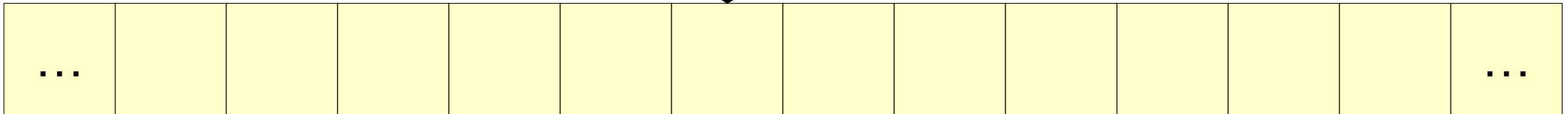
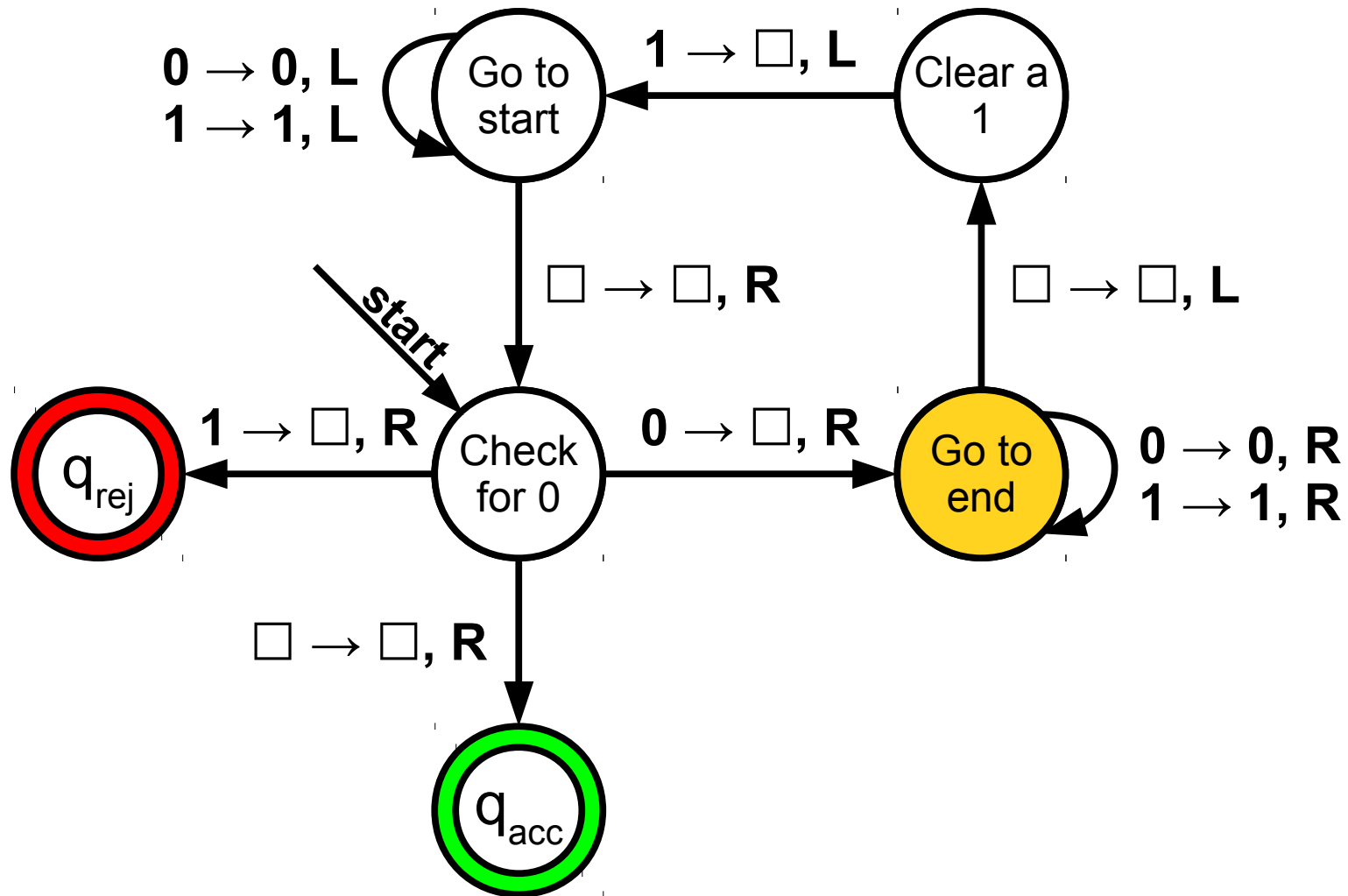


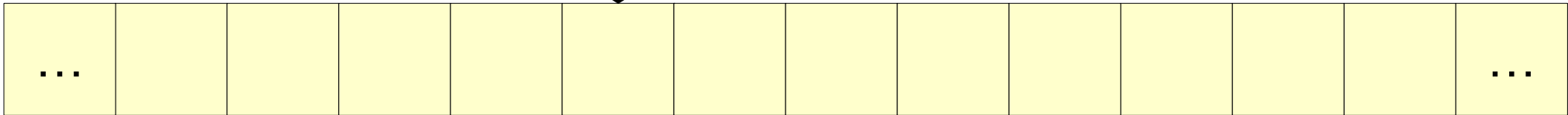
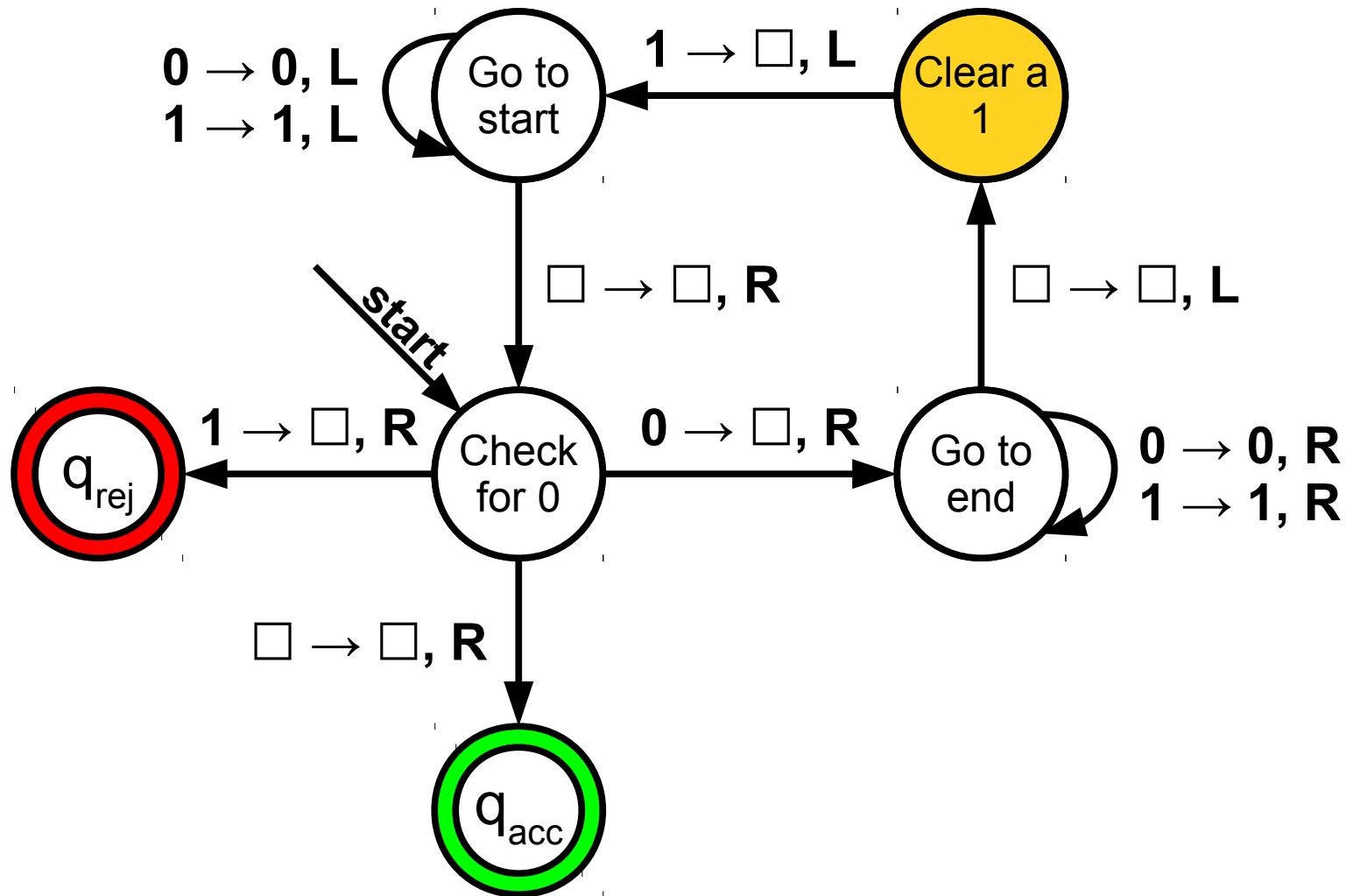






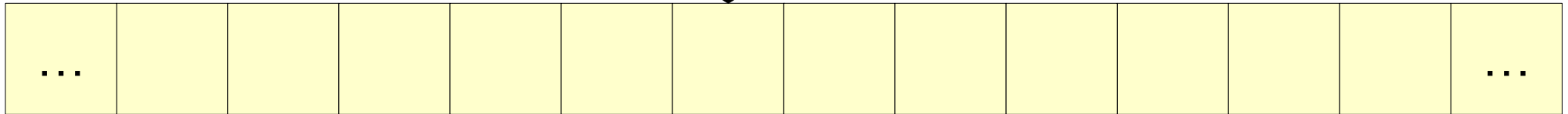
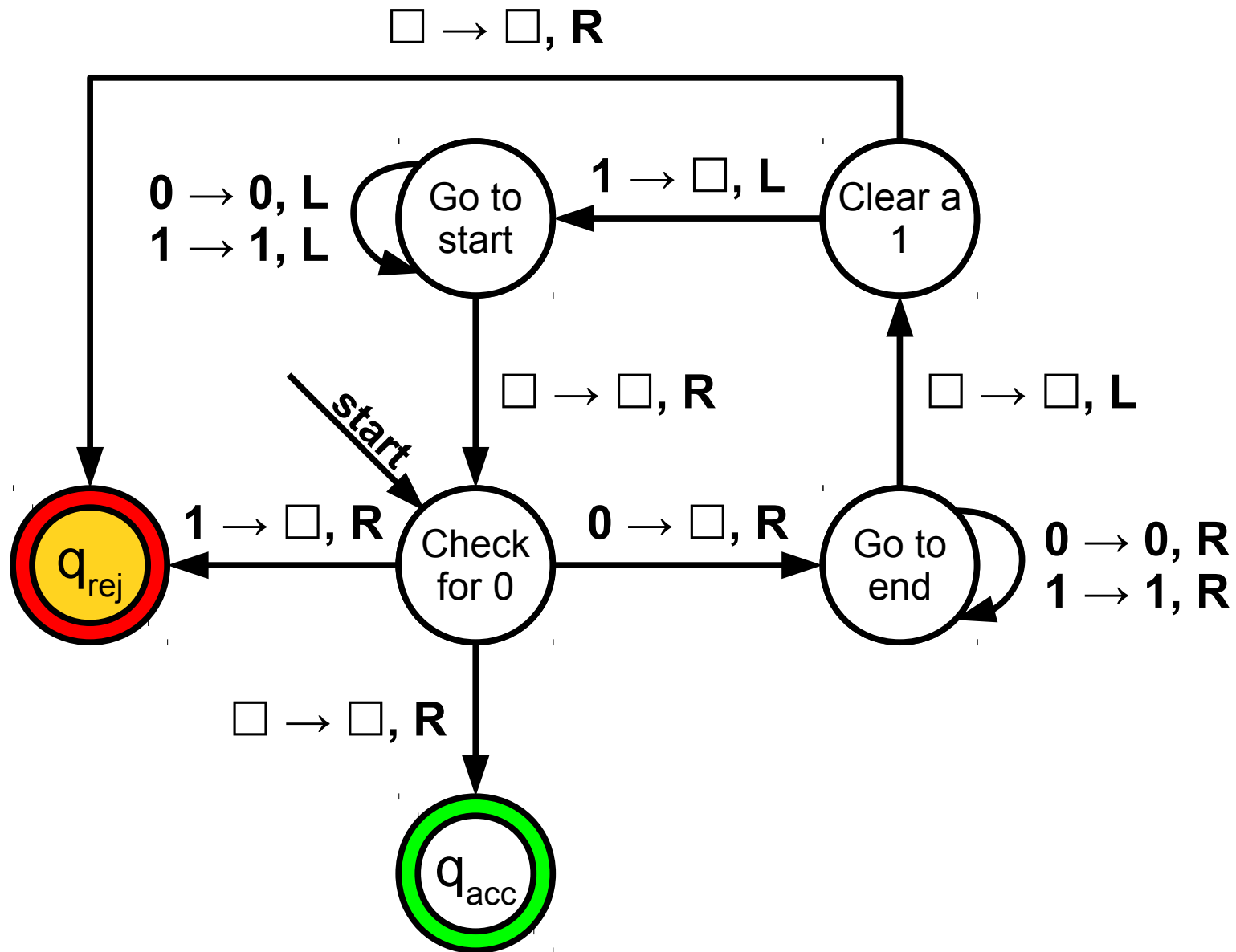




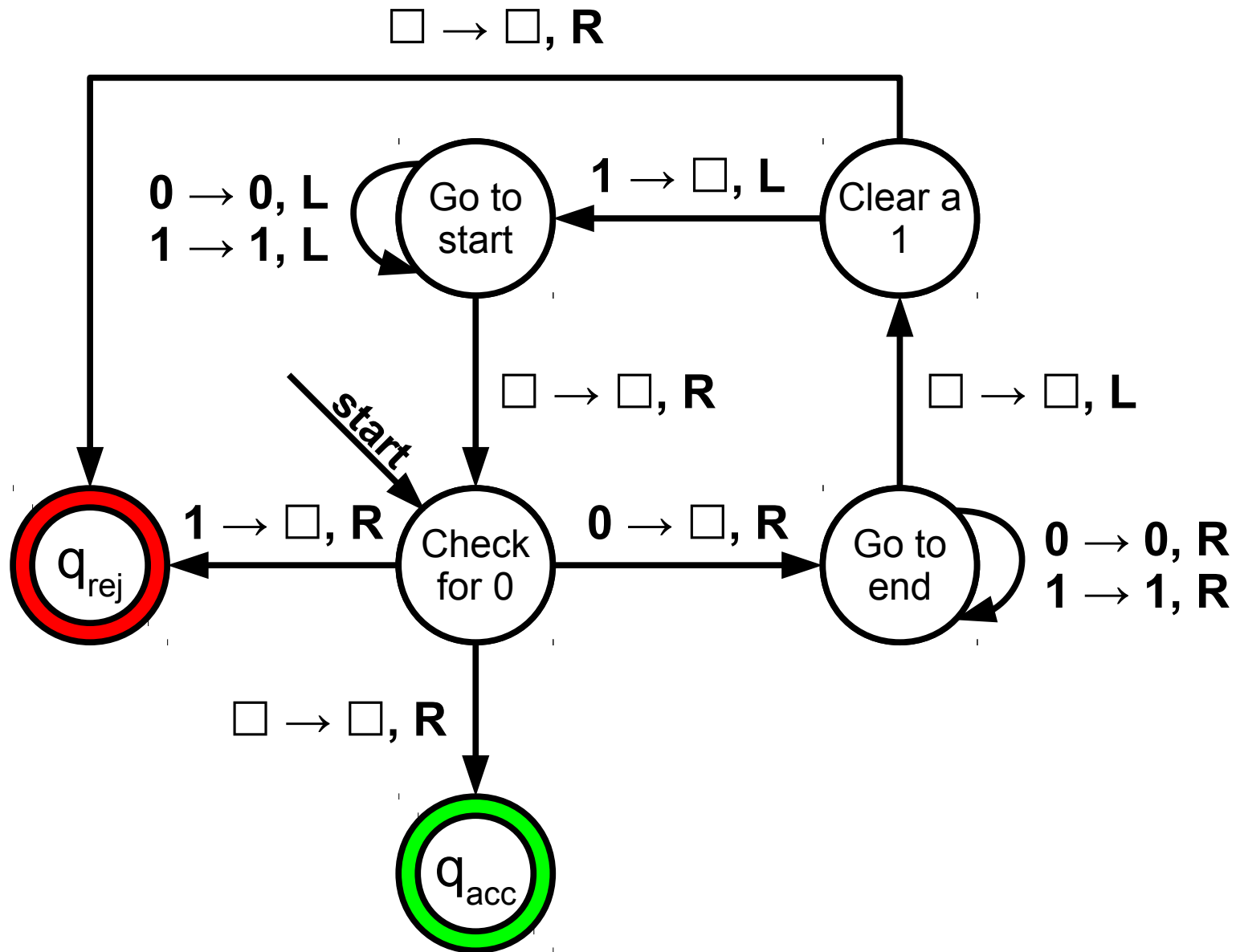


















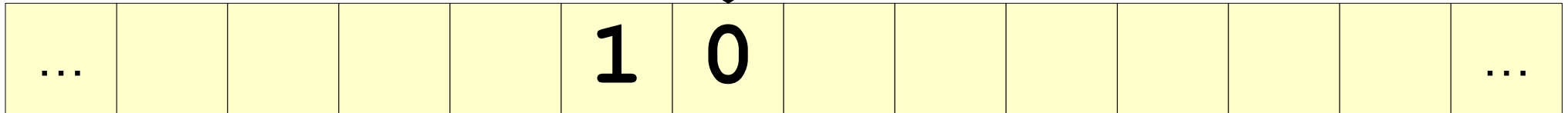
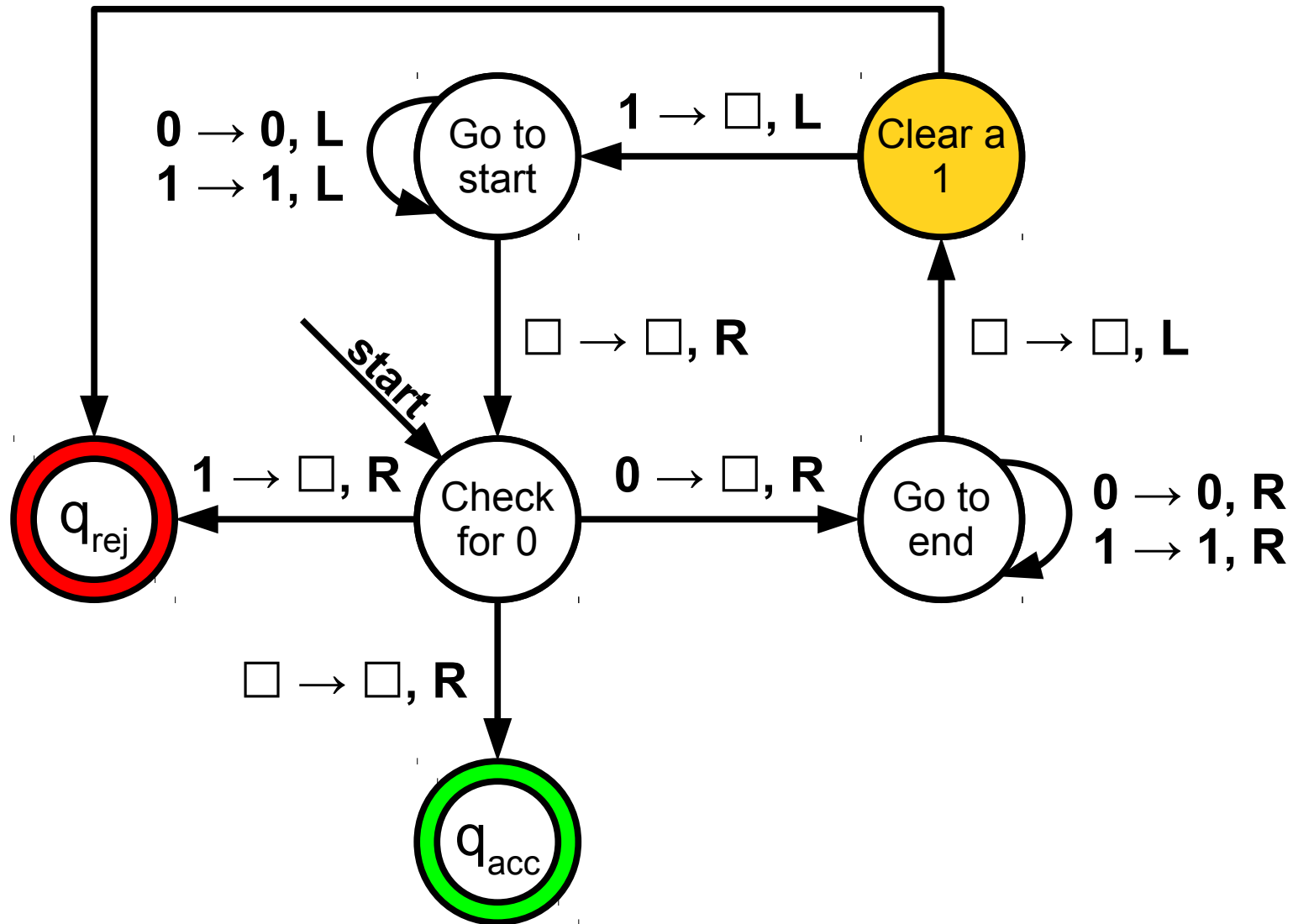


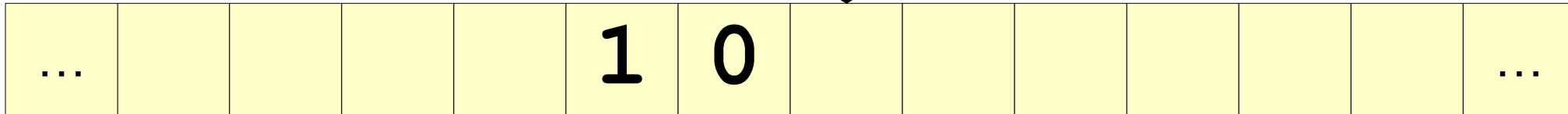
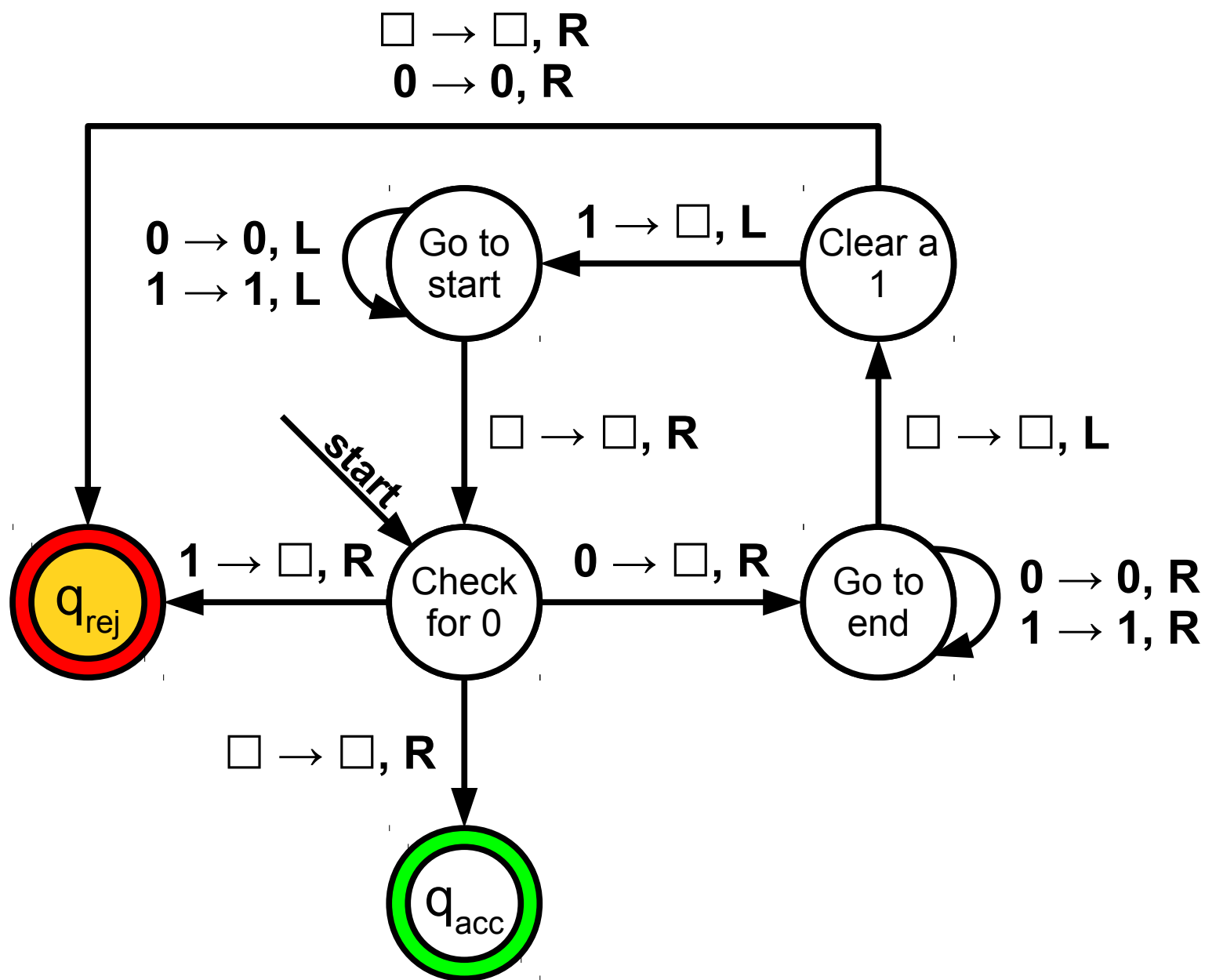


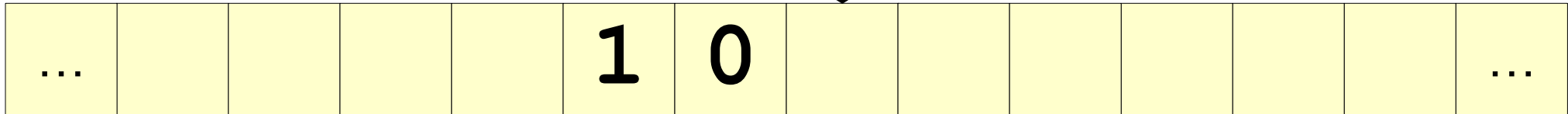
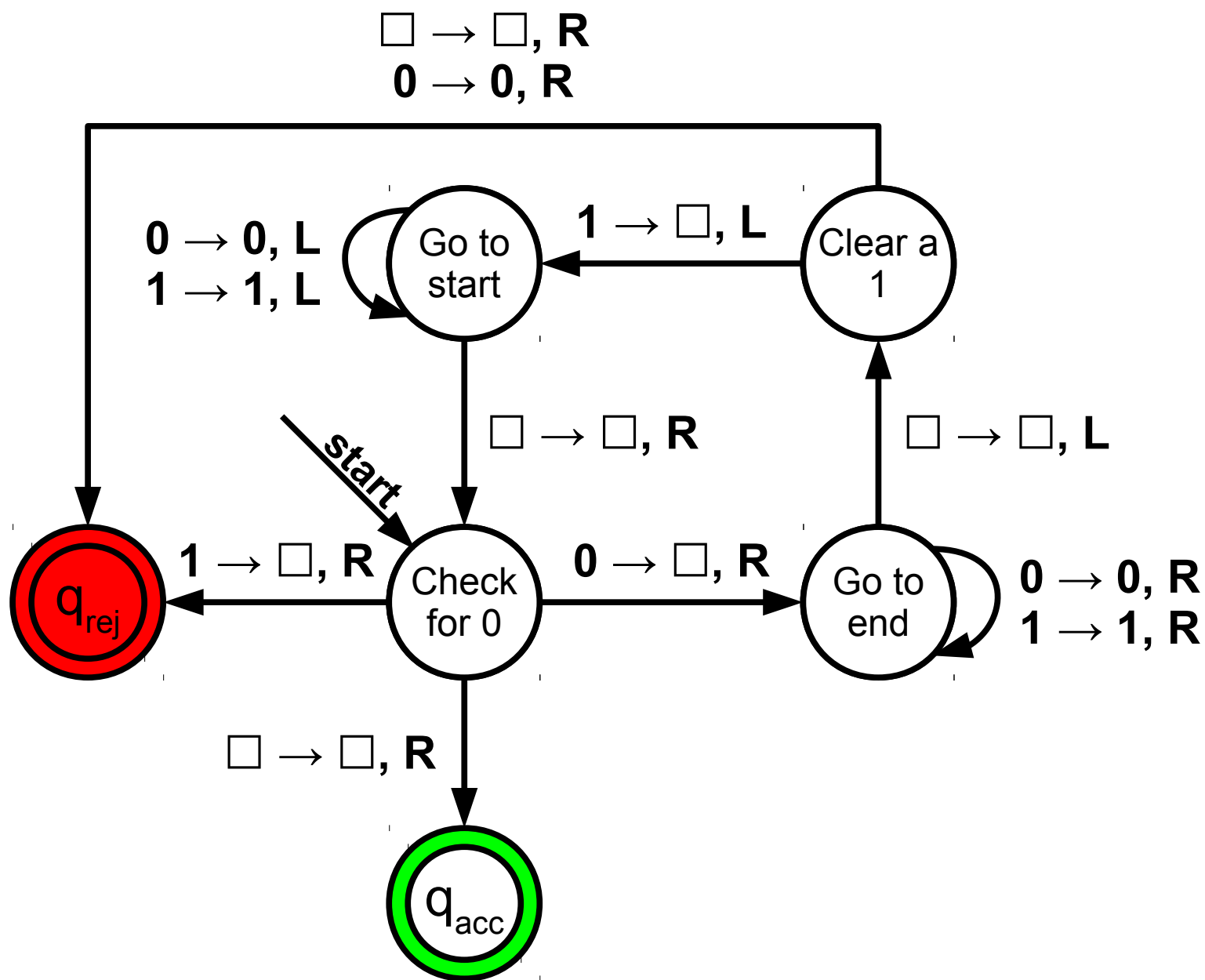


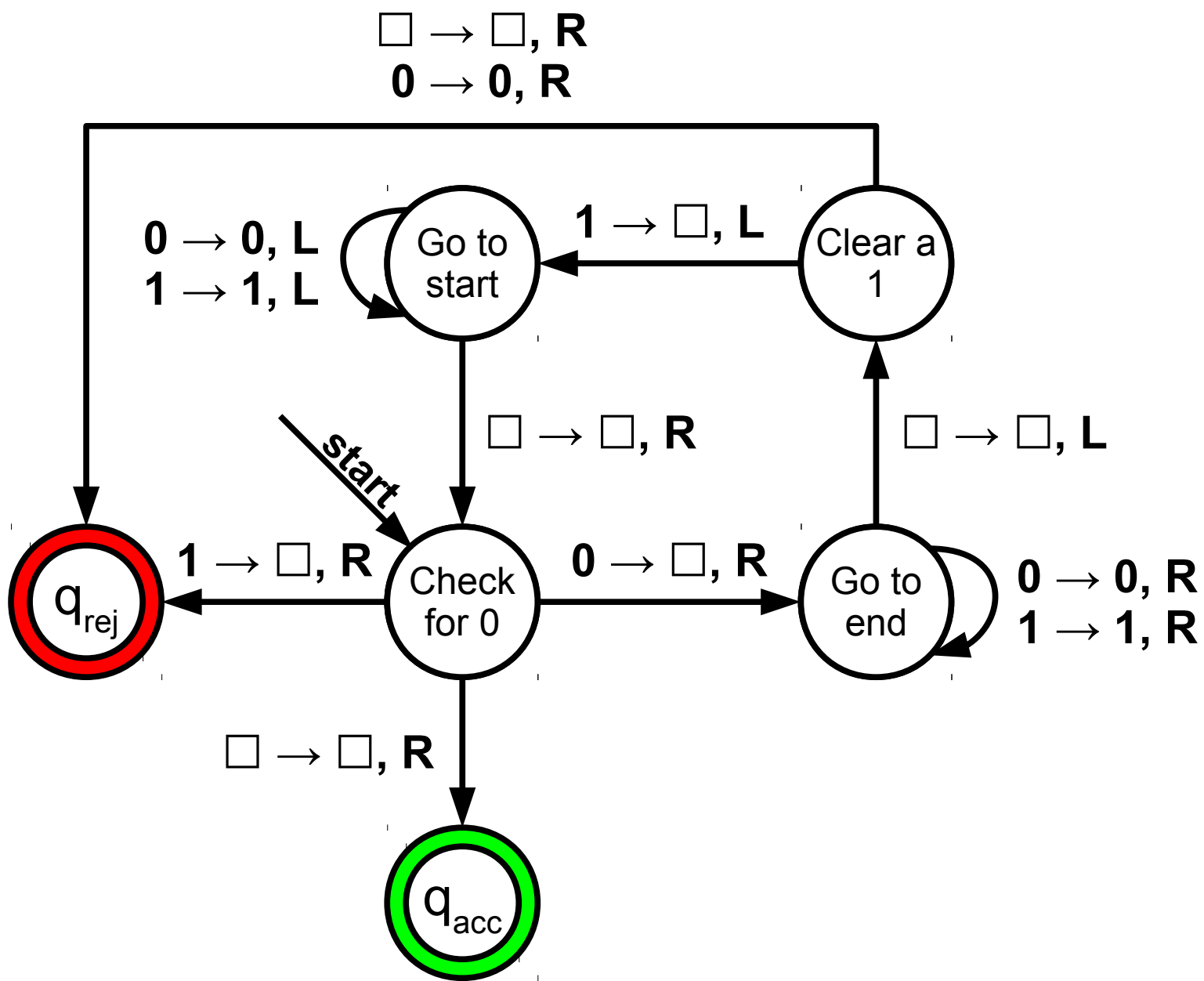


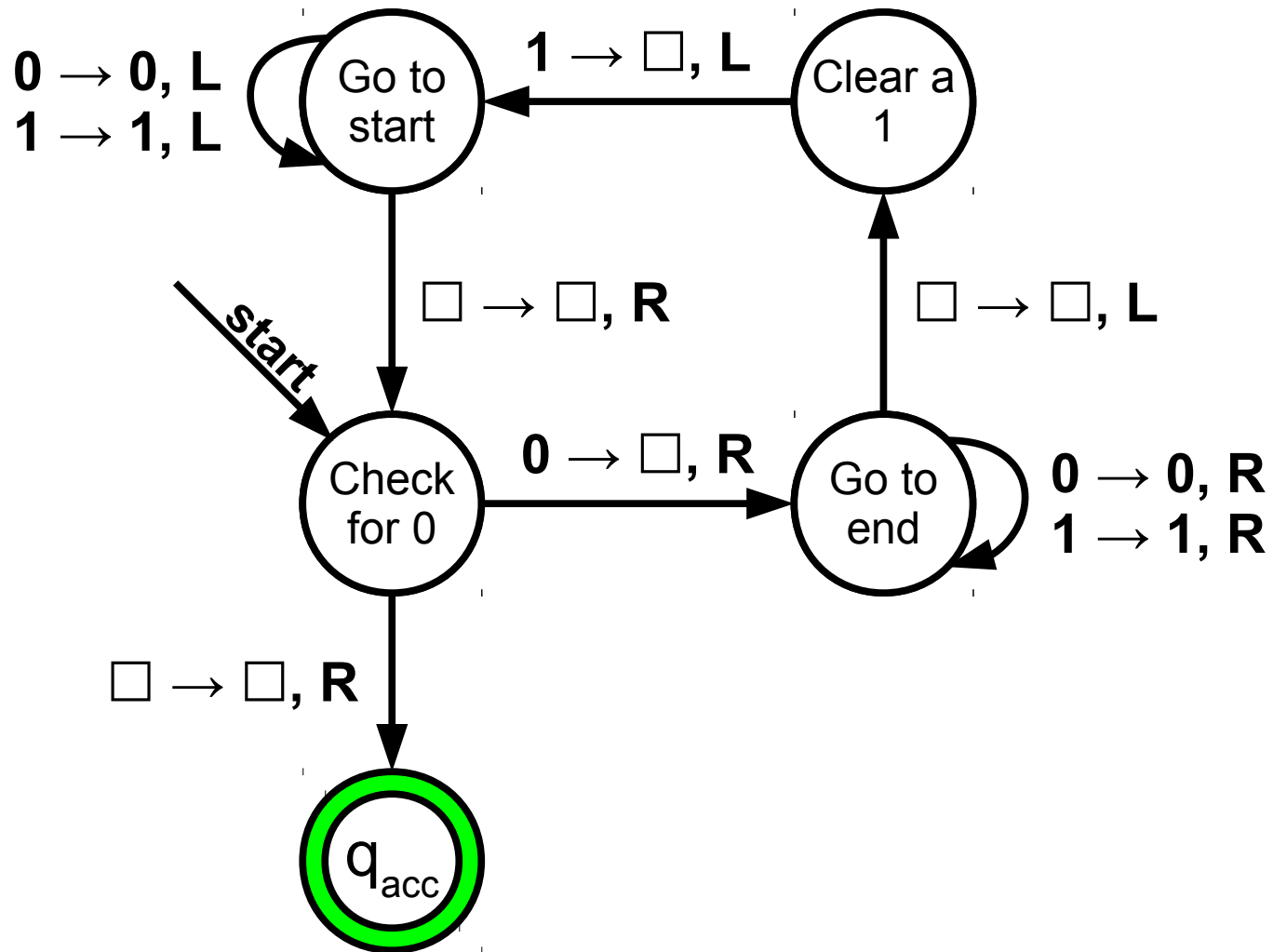
$\square \rightarrow \square, R$   
 $0 \rightarrow 0, R$









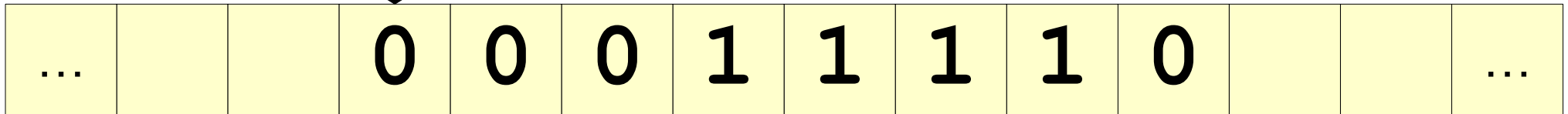


# Another TM Design

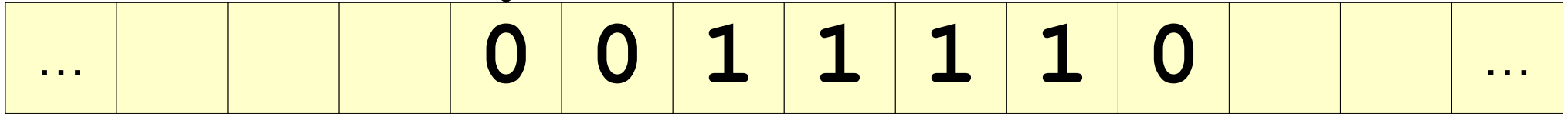
- We've designed a TM for  $\{0^n 1^n \mid n \in \mathbb{N}\}$ .
- Consider this language over  $\Sigma = \{0, 1\}$ :  
$$L = \{ w \in \Sigma^* \mid w \text{ has the same number of } 0\text{s and } 1\text{s} \}$$
- This language is also not regular, but it is context-free.
- How might we design a TM for it?



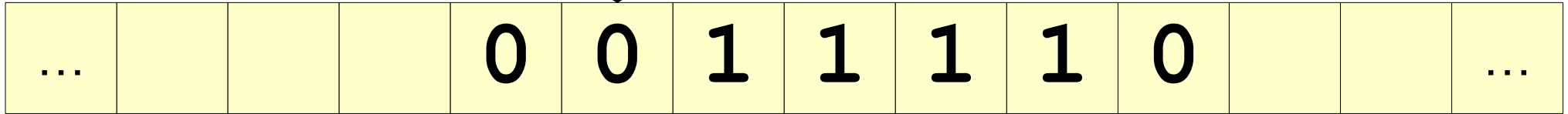
# A Caveat



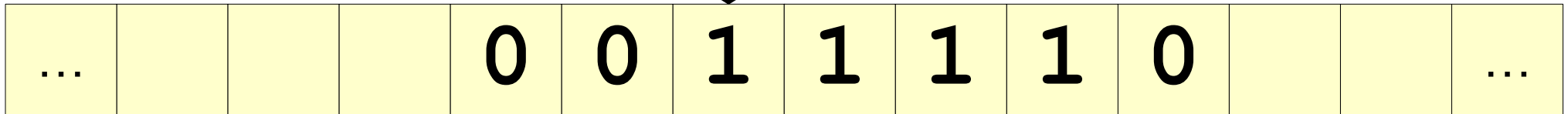
# A Caveat



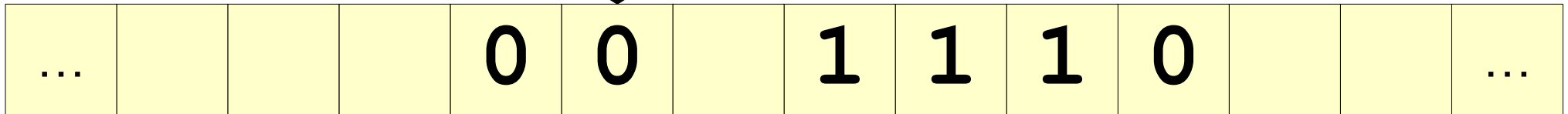
# A Caveat



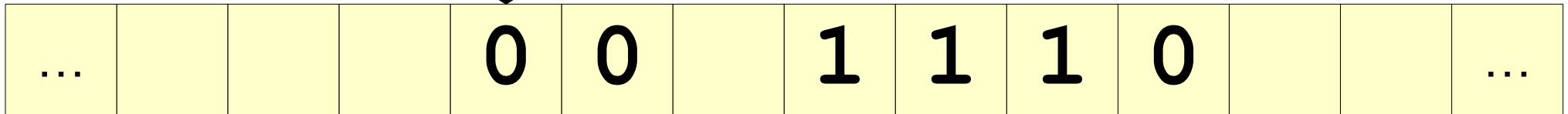
# A Caveat



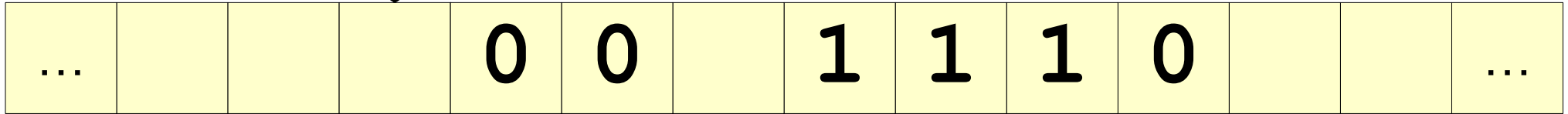
# A Caveat



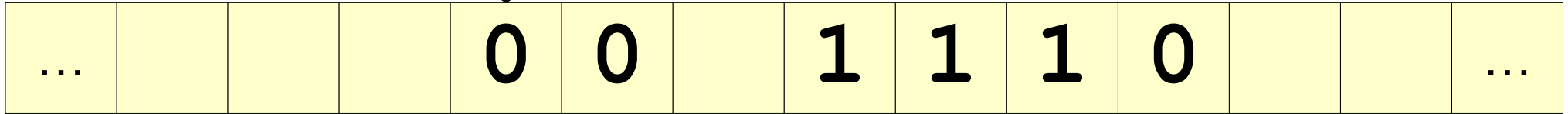
# A Caveat



# A Caveat

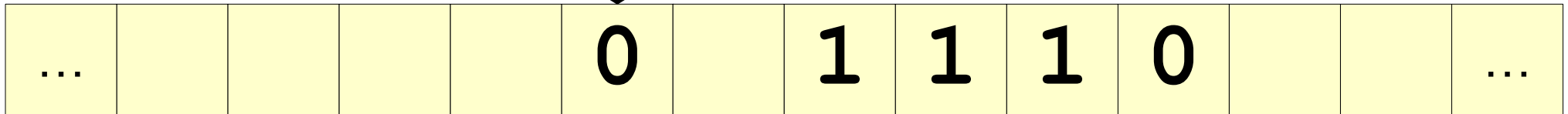


# A Caveat

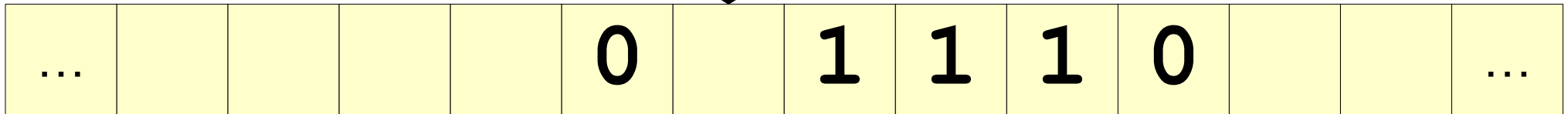




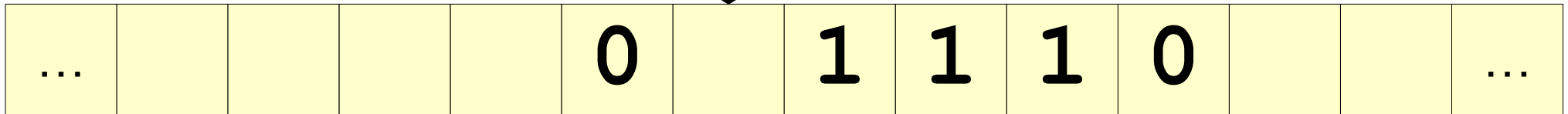
# A Caveat



# A Caveat

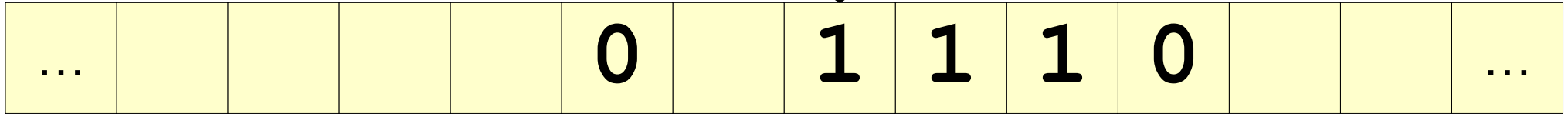


# A Caveat

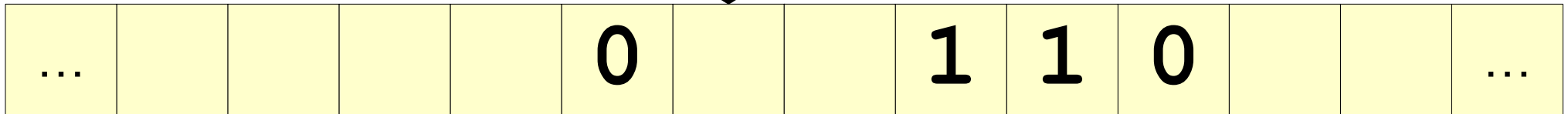


How do we know that  
this blank isn't one of  
the infinitely many  
blanks after our input  
string?

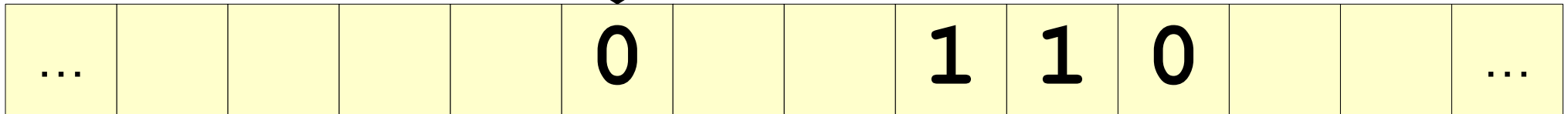
# A Caveat



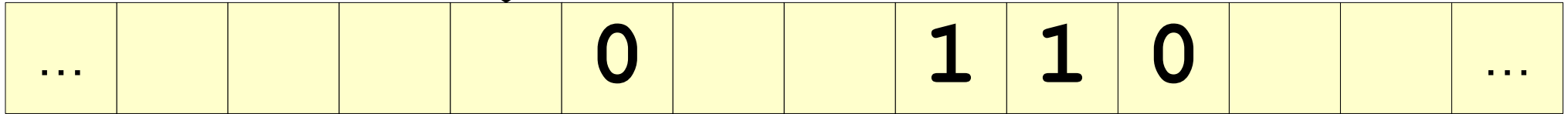
# A Caveat



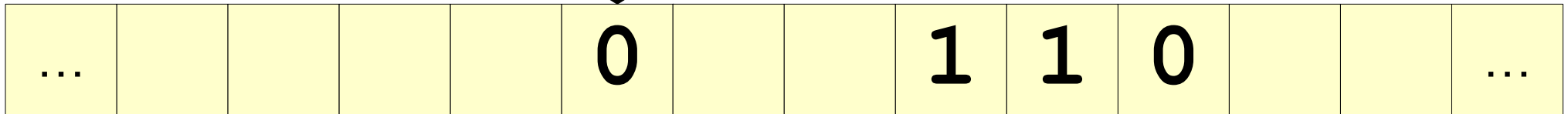
# A Caveat



# A Caveat

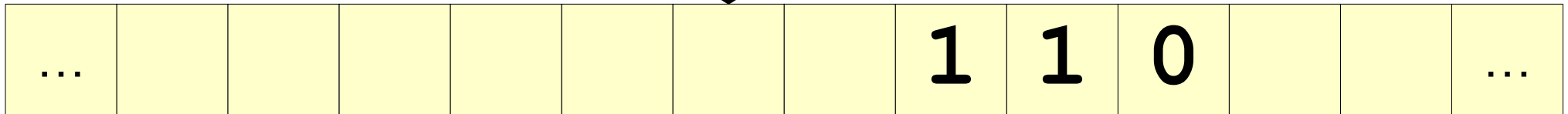


# A Caveat

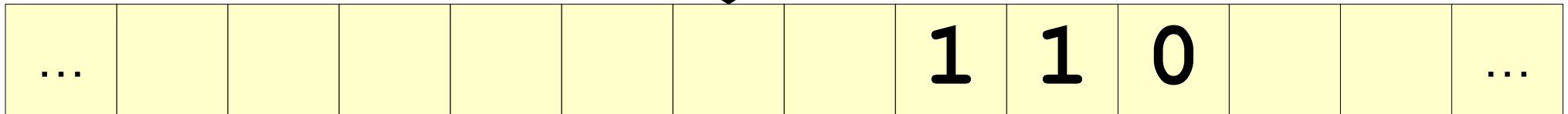




# A Caveat

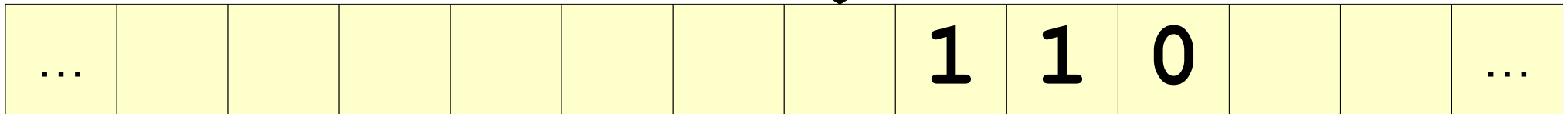


# A Caveat

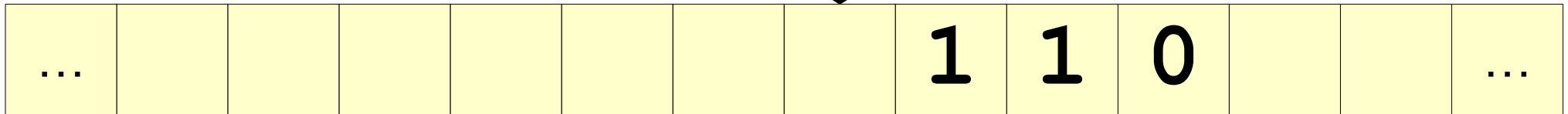


How do we know that this blank isn't one of the infinitely many blanks after our input string?

# A Caveat

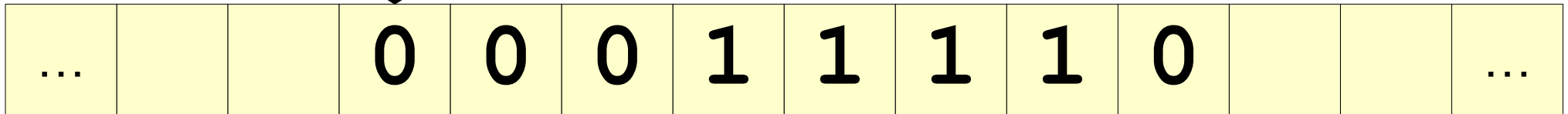


# A Caveat



How do we know that  
this blank isn't one of  
the infinitely many  
blanks after our input  
string?

# One Solution



# One Solution



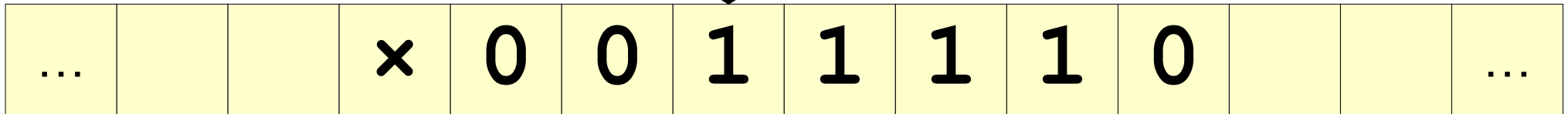
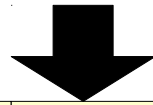
...			x	0	0	1	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

# One Solution



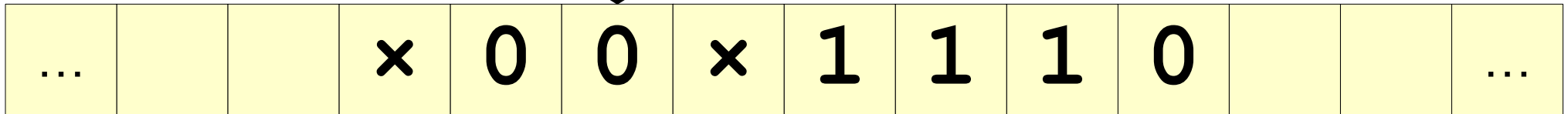
...			x	0	0	1	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

# One Solution

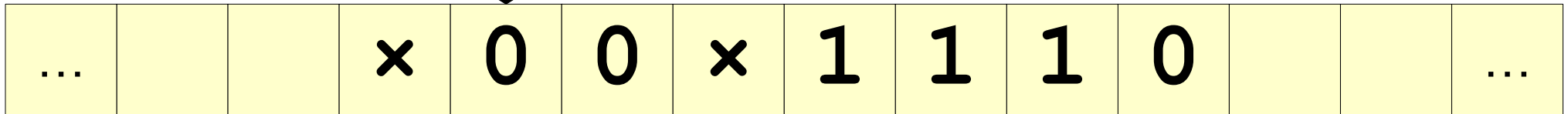




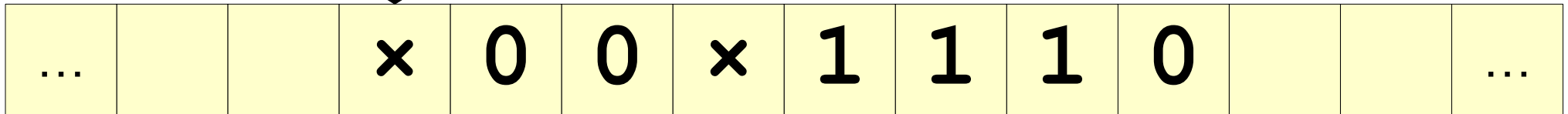
# One Solution



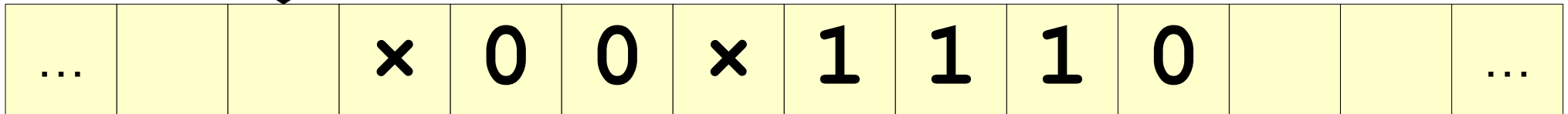
# One Solution



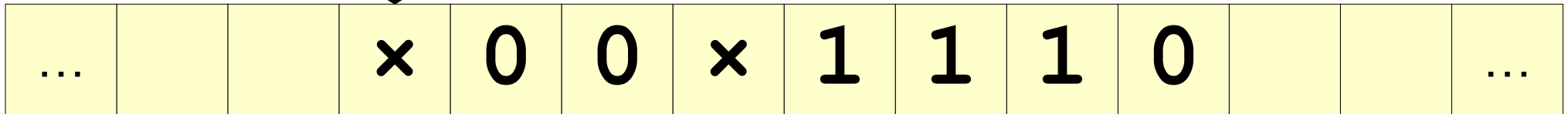
# One Solution



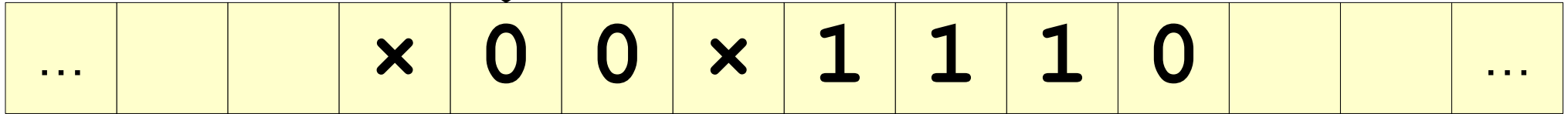
# One Solution



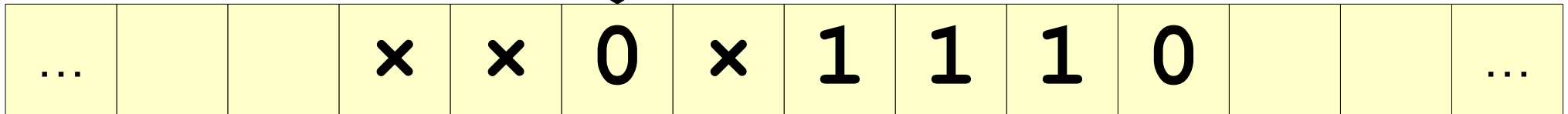
# One Solution



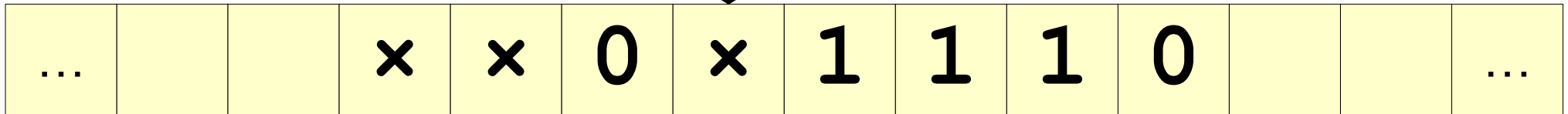
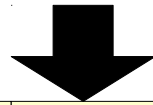
# One Solution



# One Solution



# One Solution



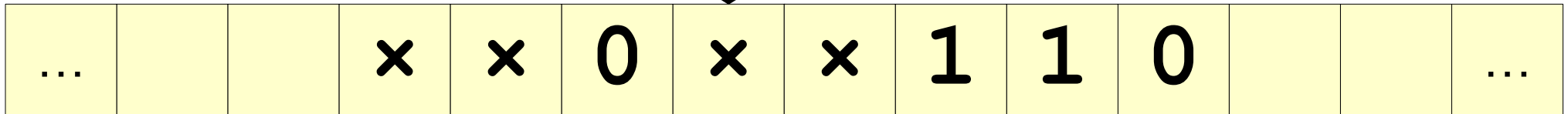


# One Solution

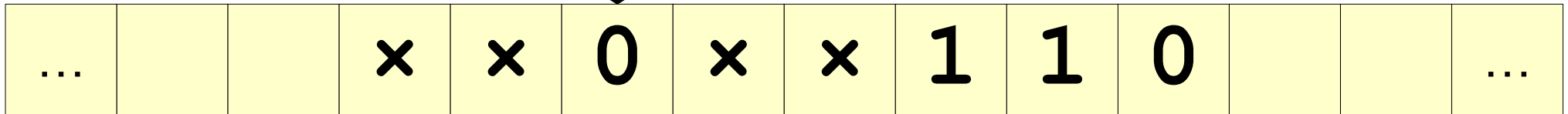


...			x	x	0	x	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

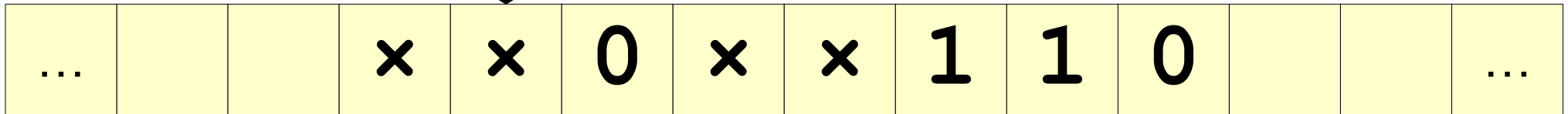
# One Solution



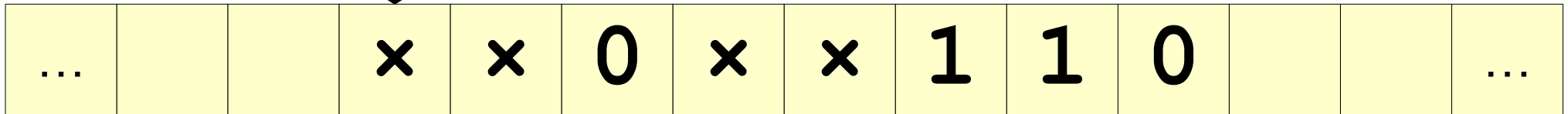
# One Solution



# One Solution



# One Solution

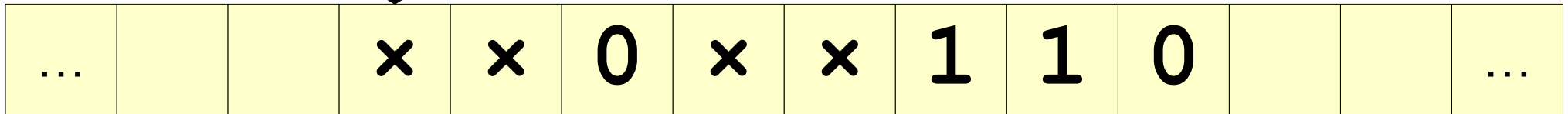


# One Solution

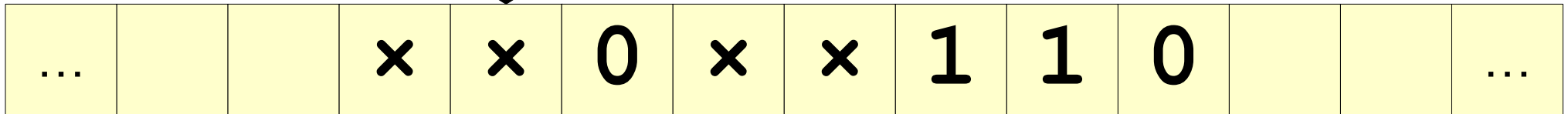


...			x	x	0	x	x	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

# One Solution

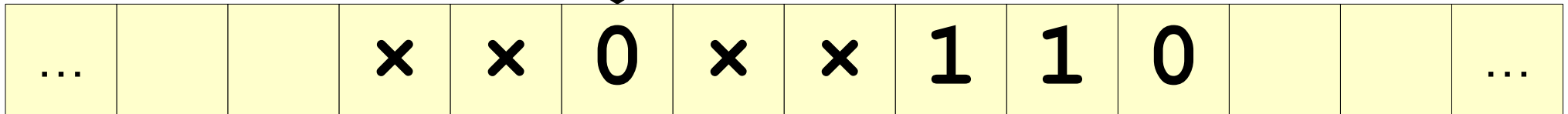


# One Solution

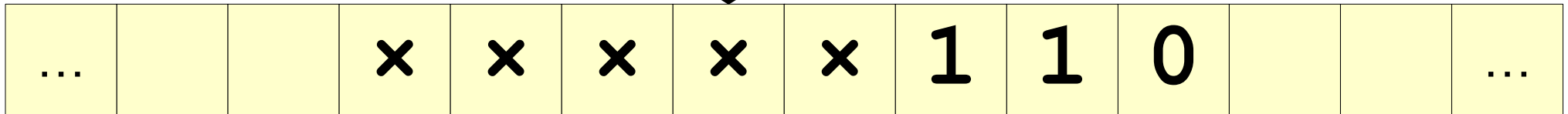




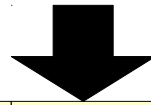
# One Solution



# One Solution

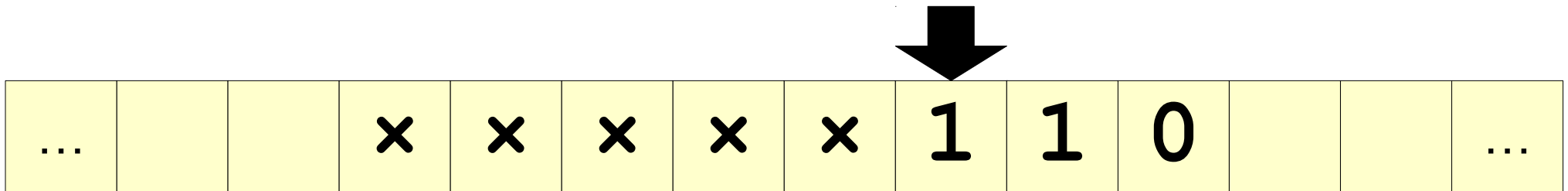


# One Solution

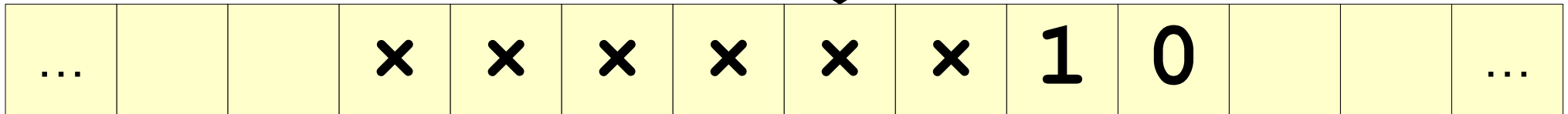


...			x	x	x	x	x	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

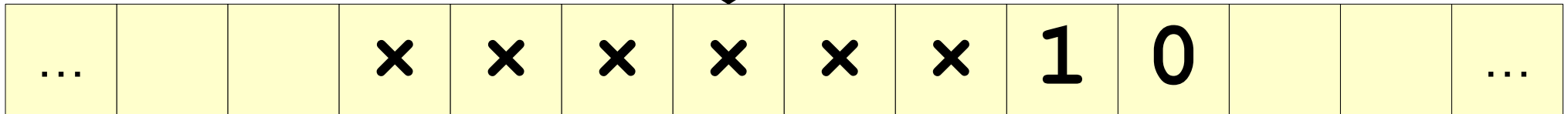
# One Solution

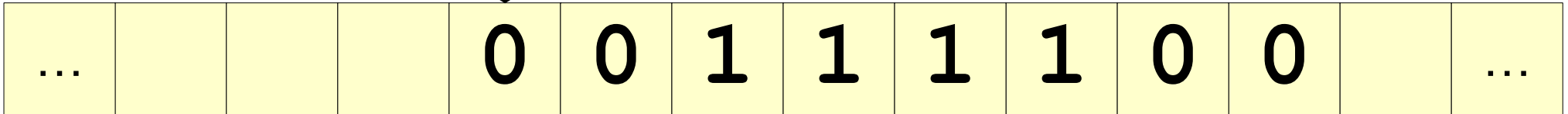
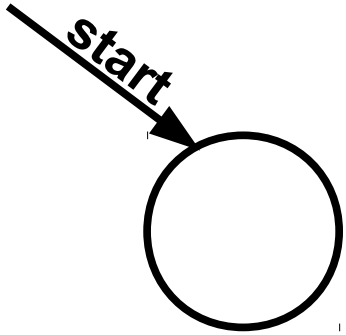


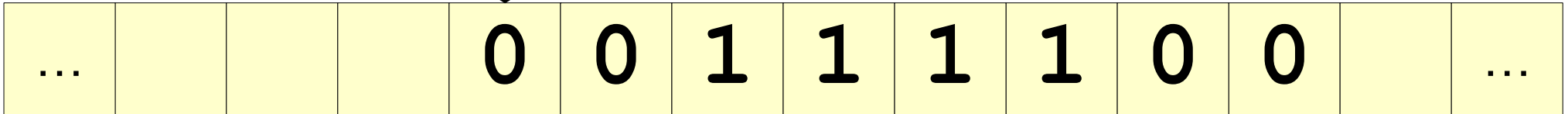
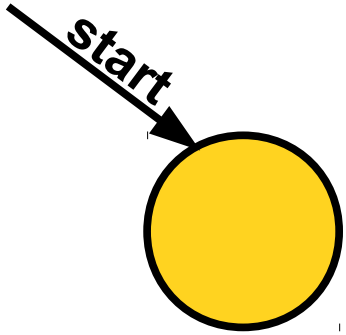
# One Solution



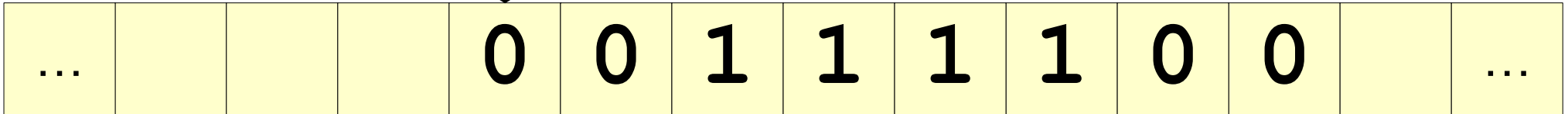
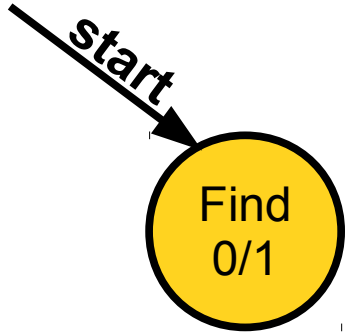
# One Solution

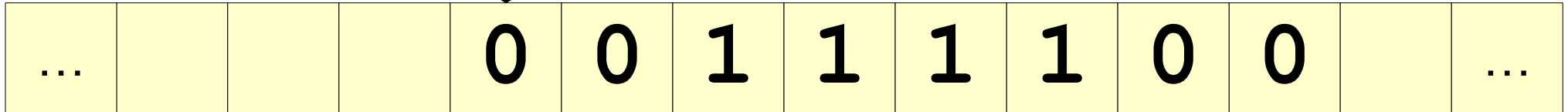
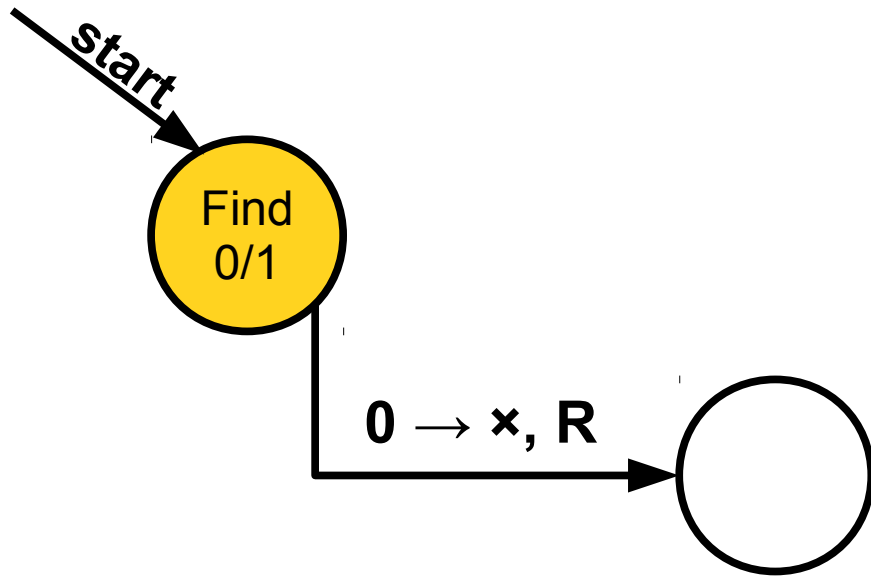


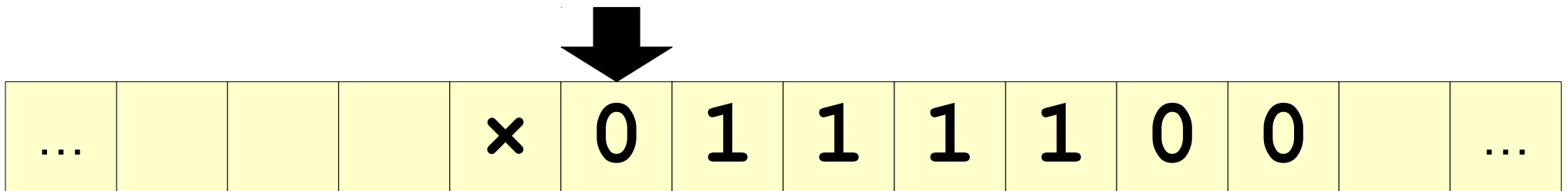
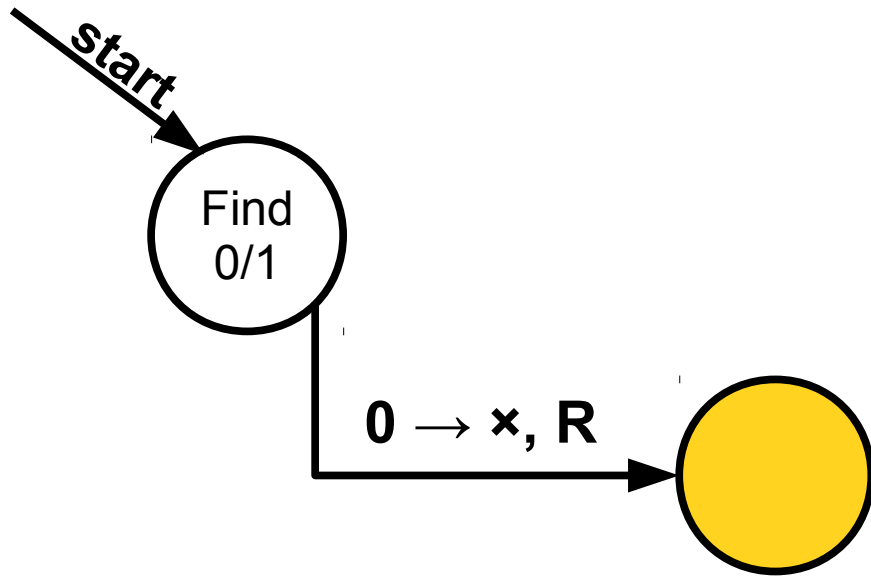


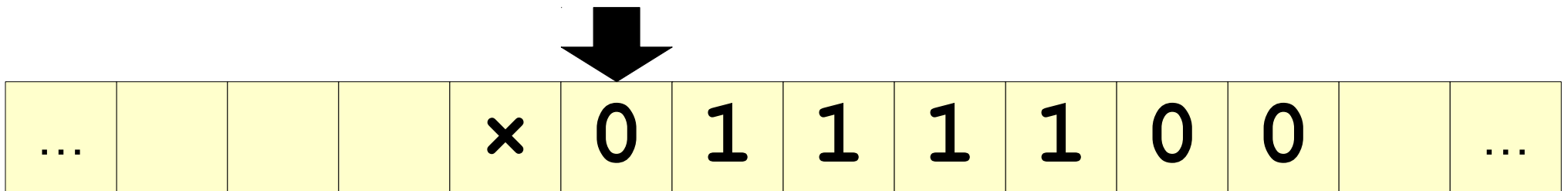
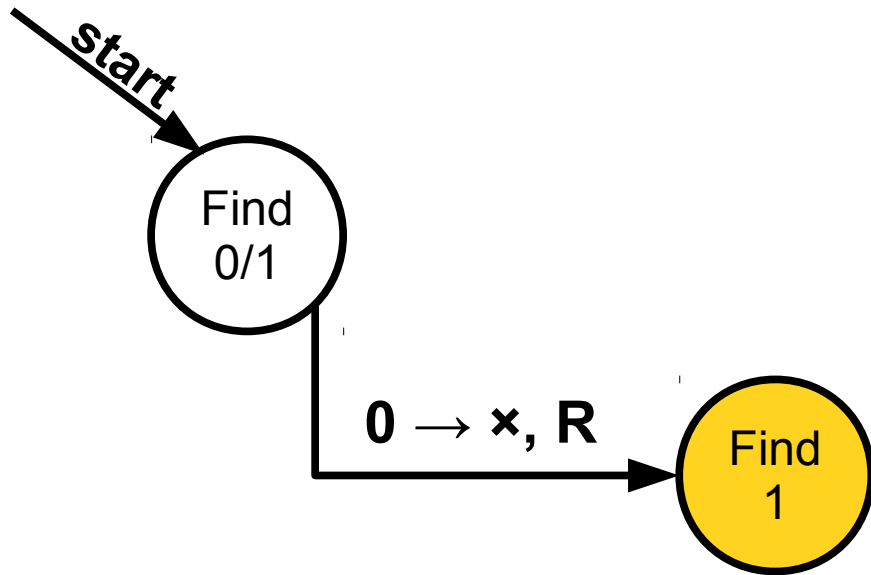


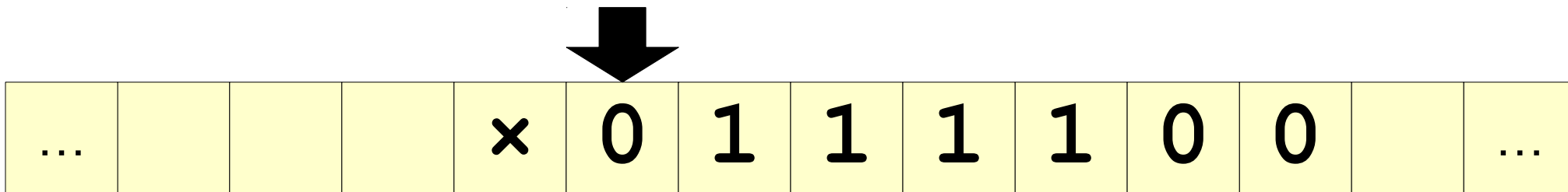
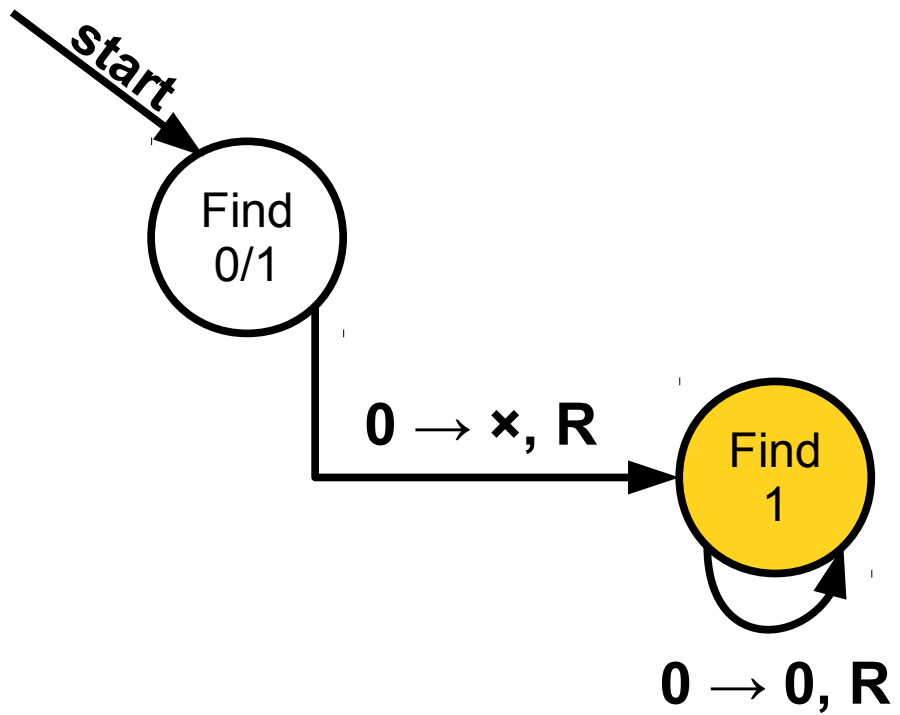


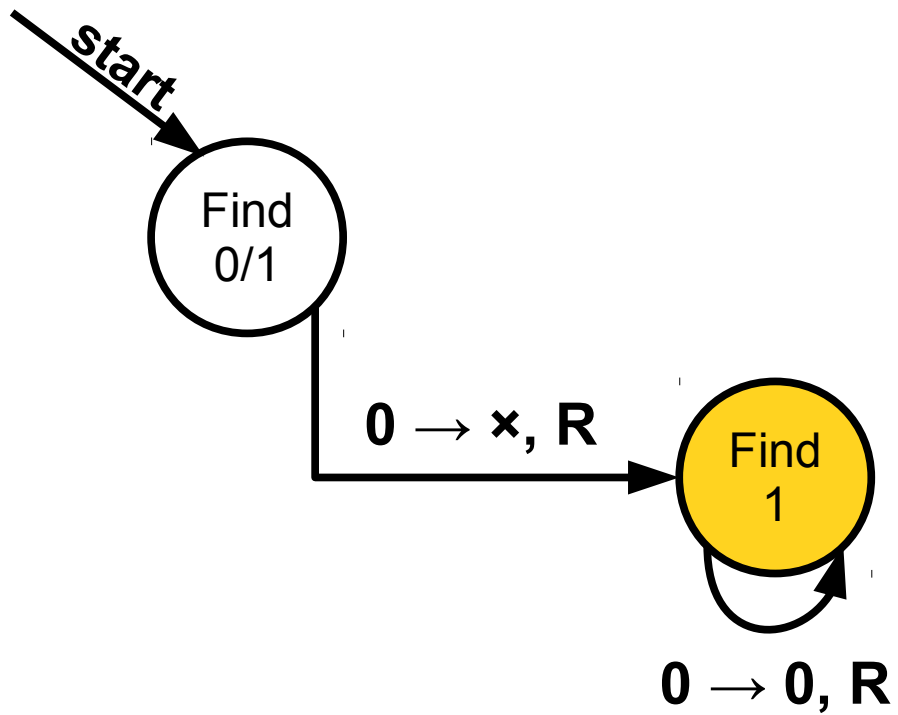


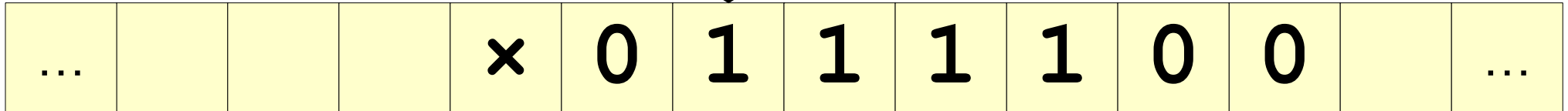
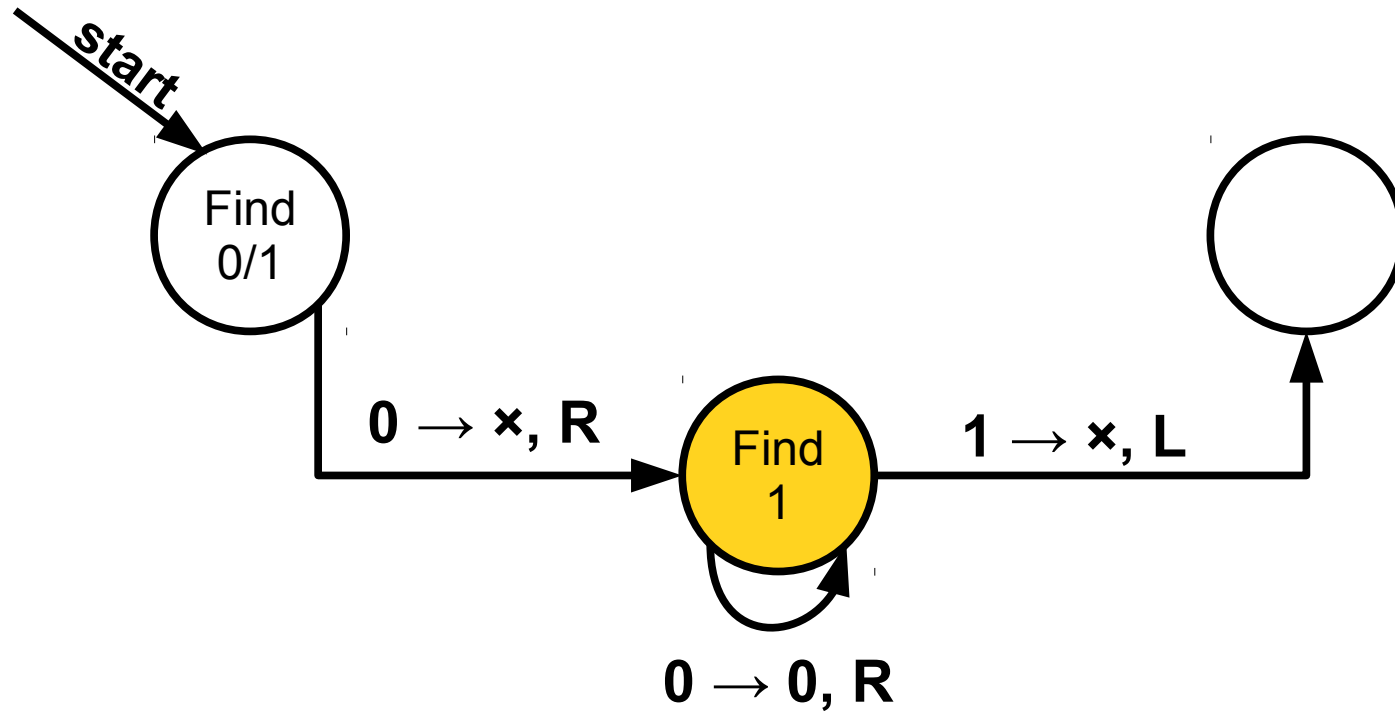


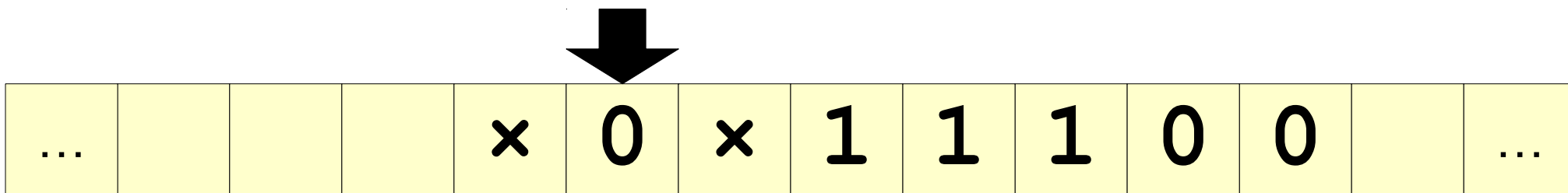
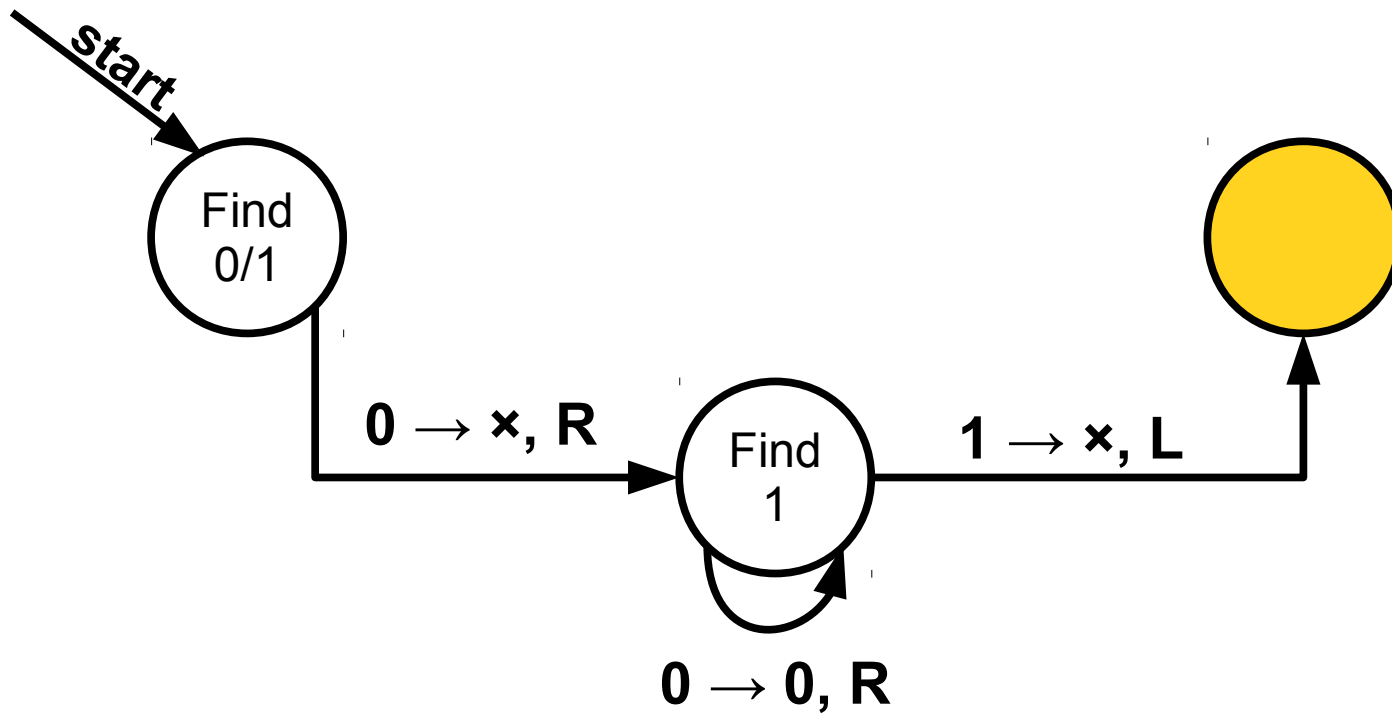




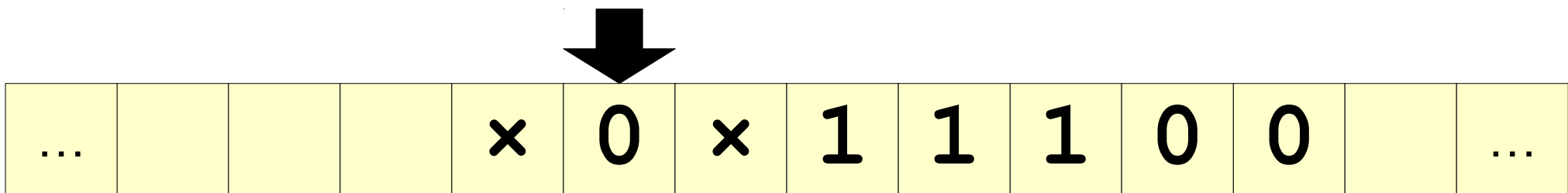
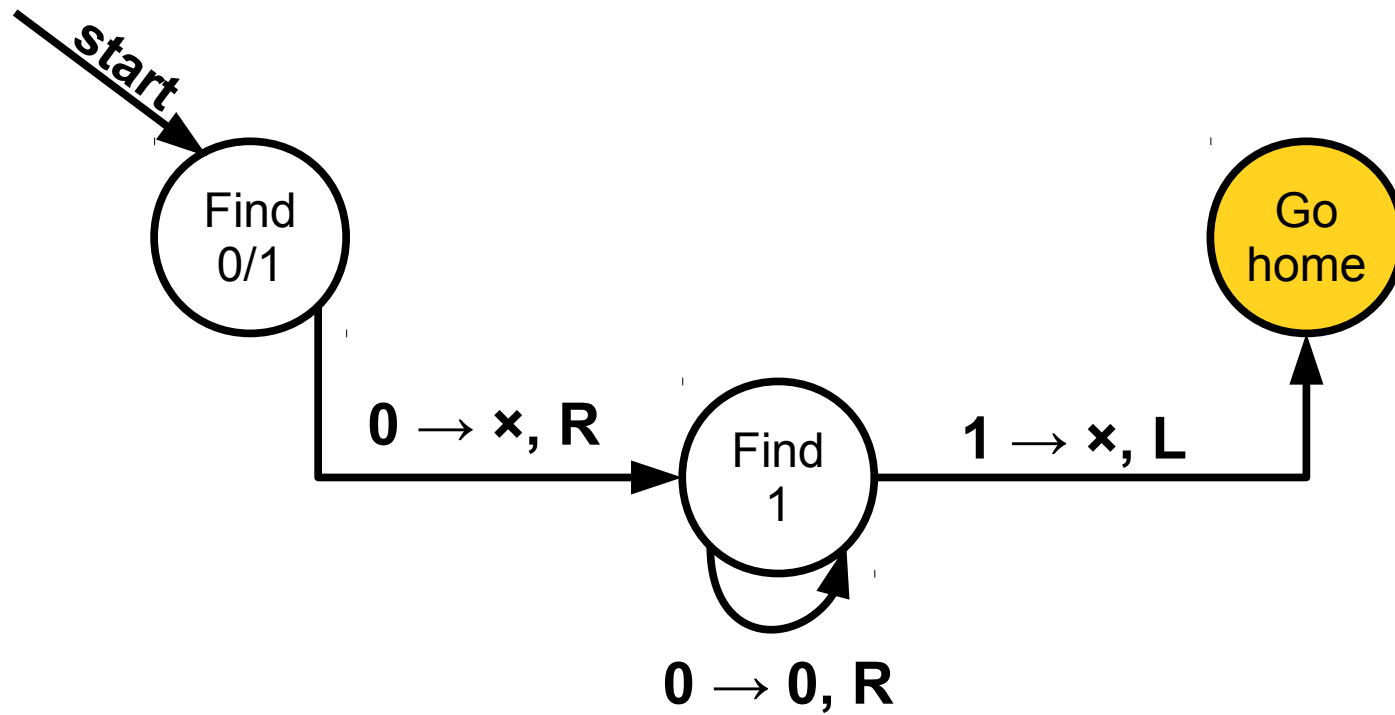


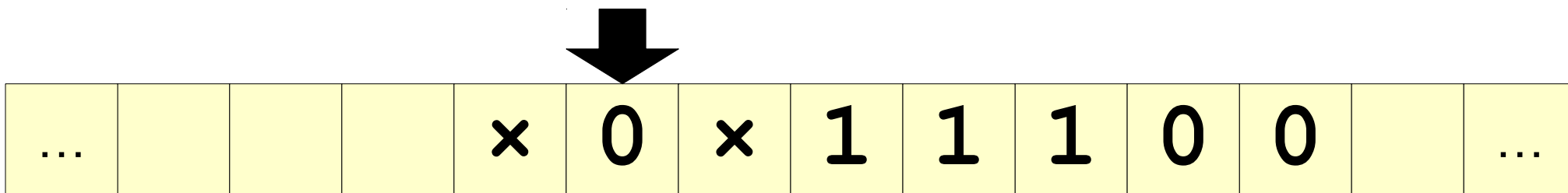
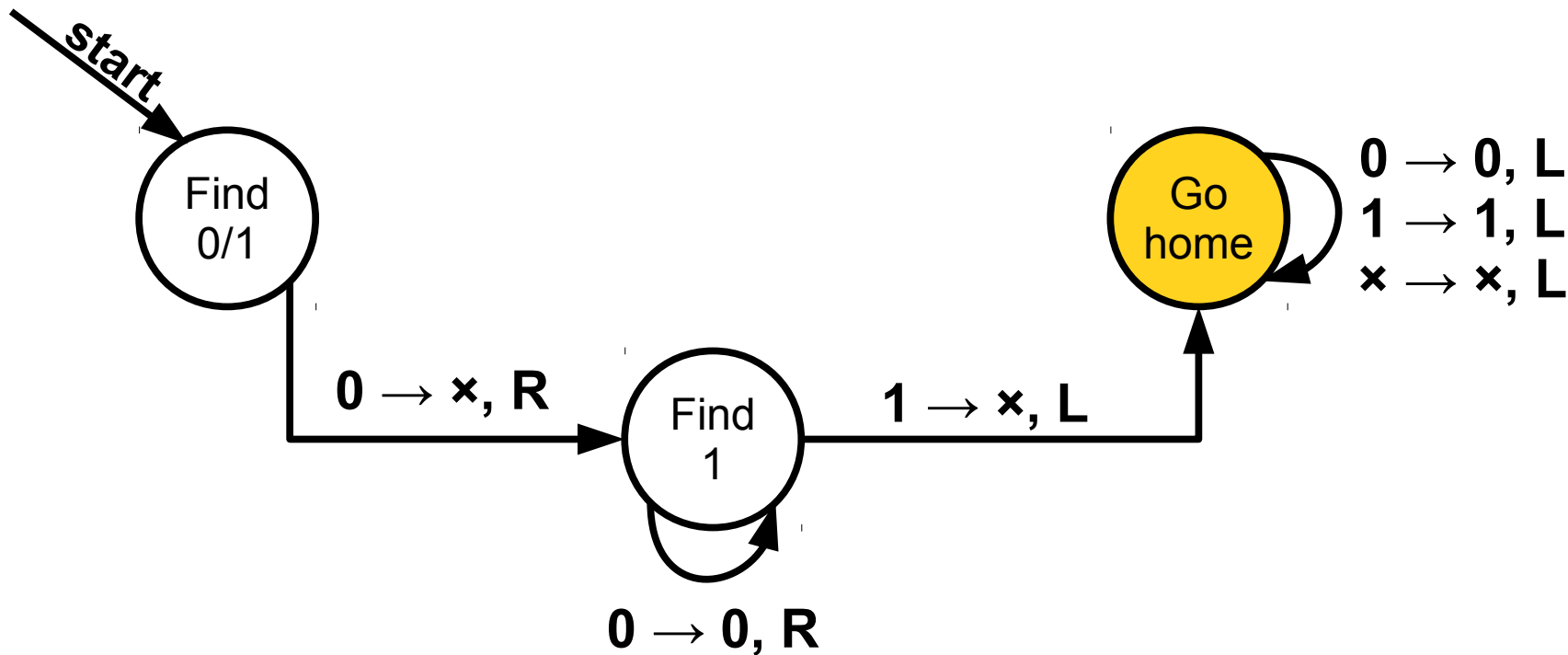


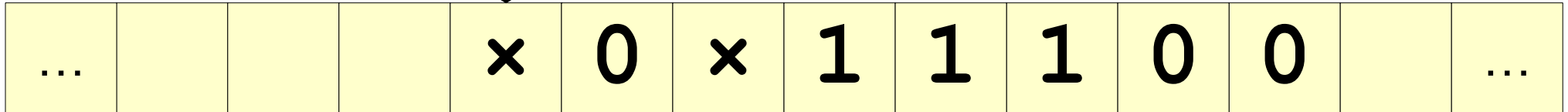
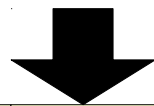
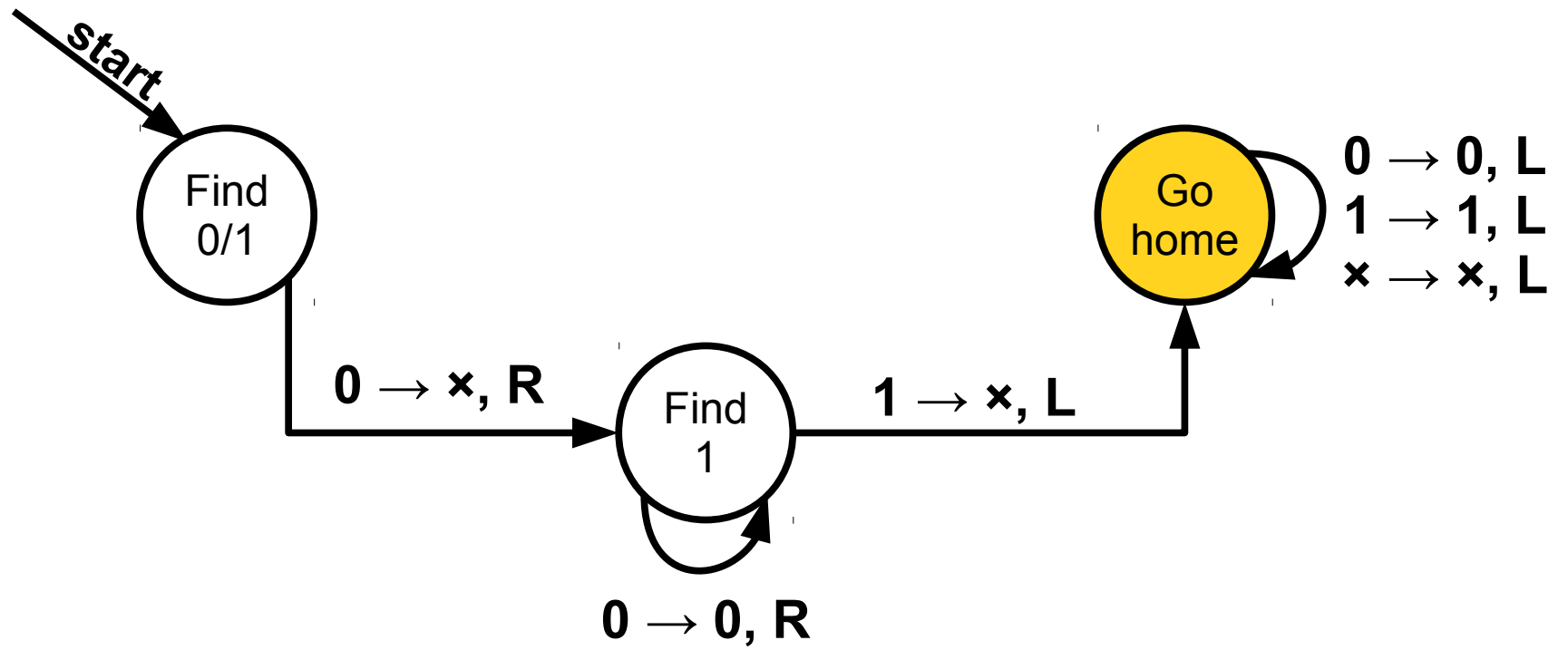


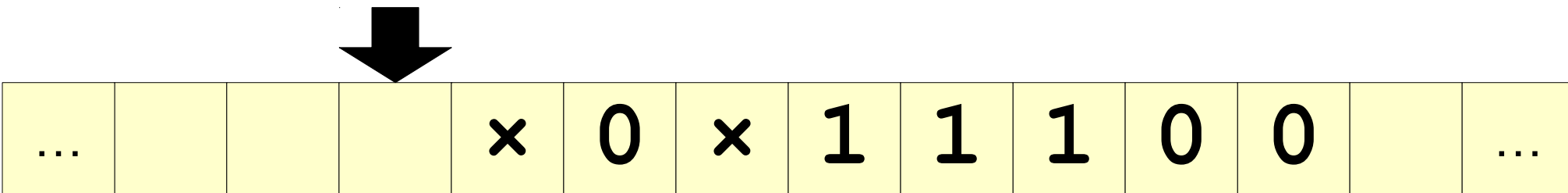
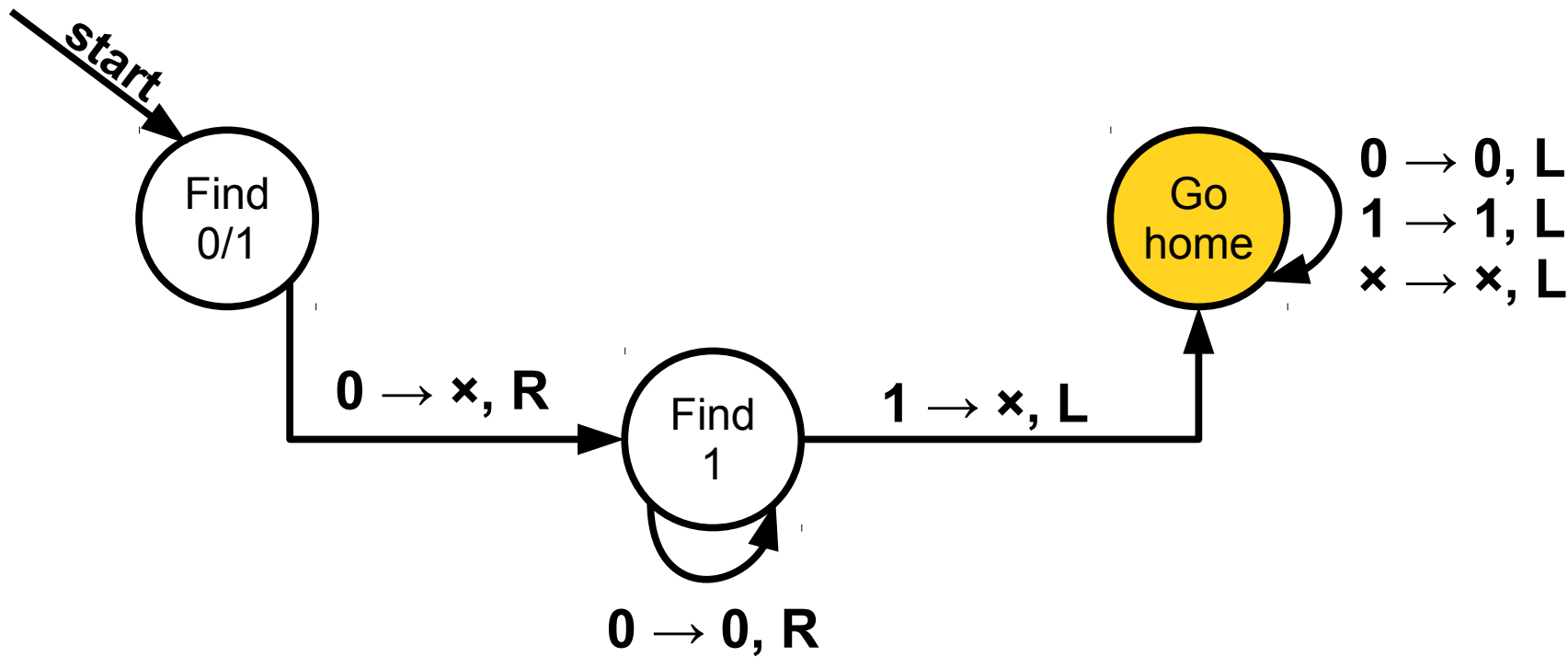


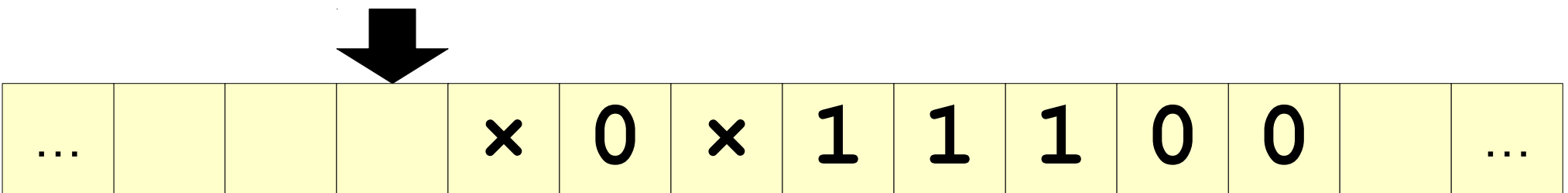
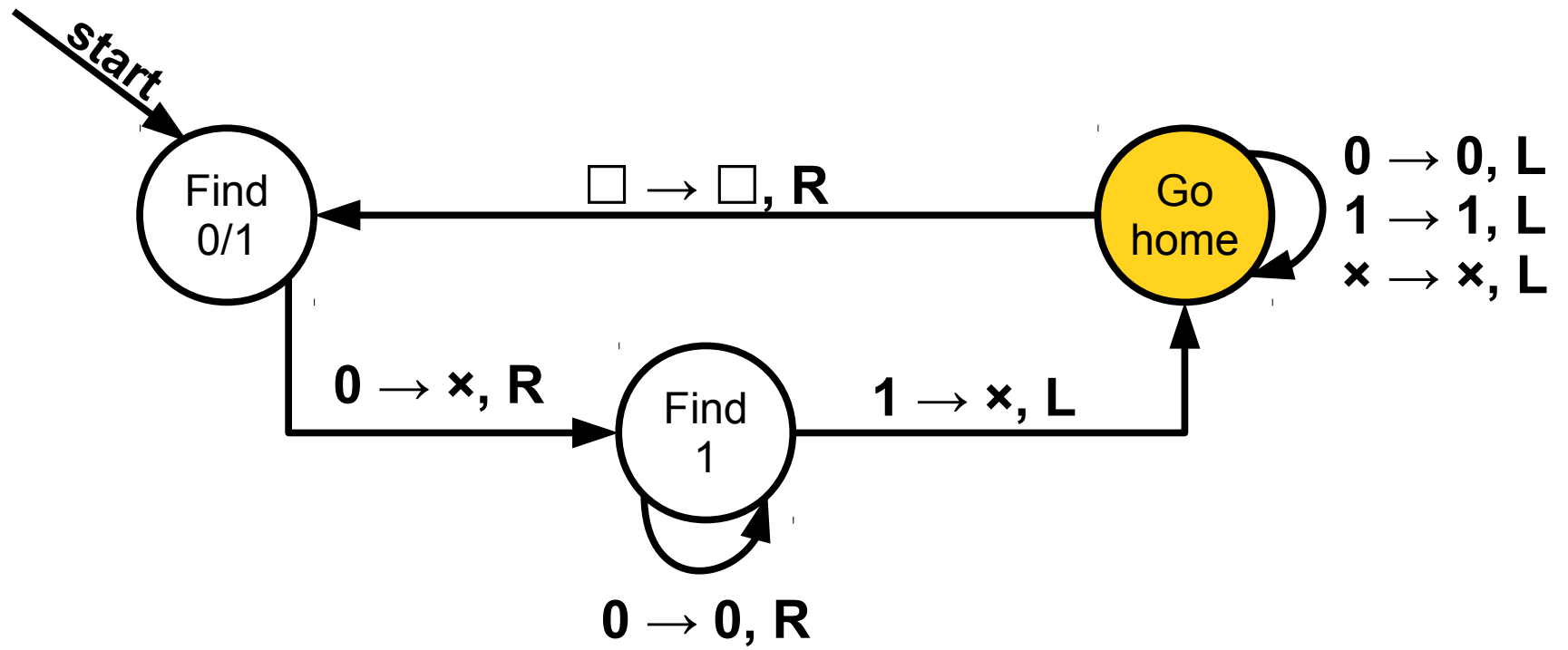


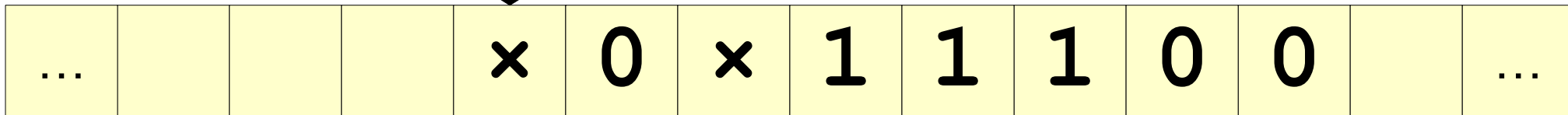
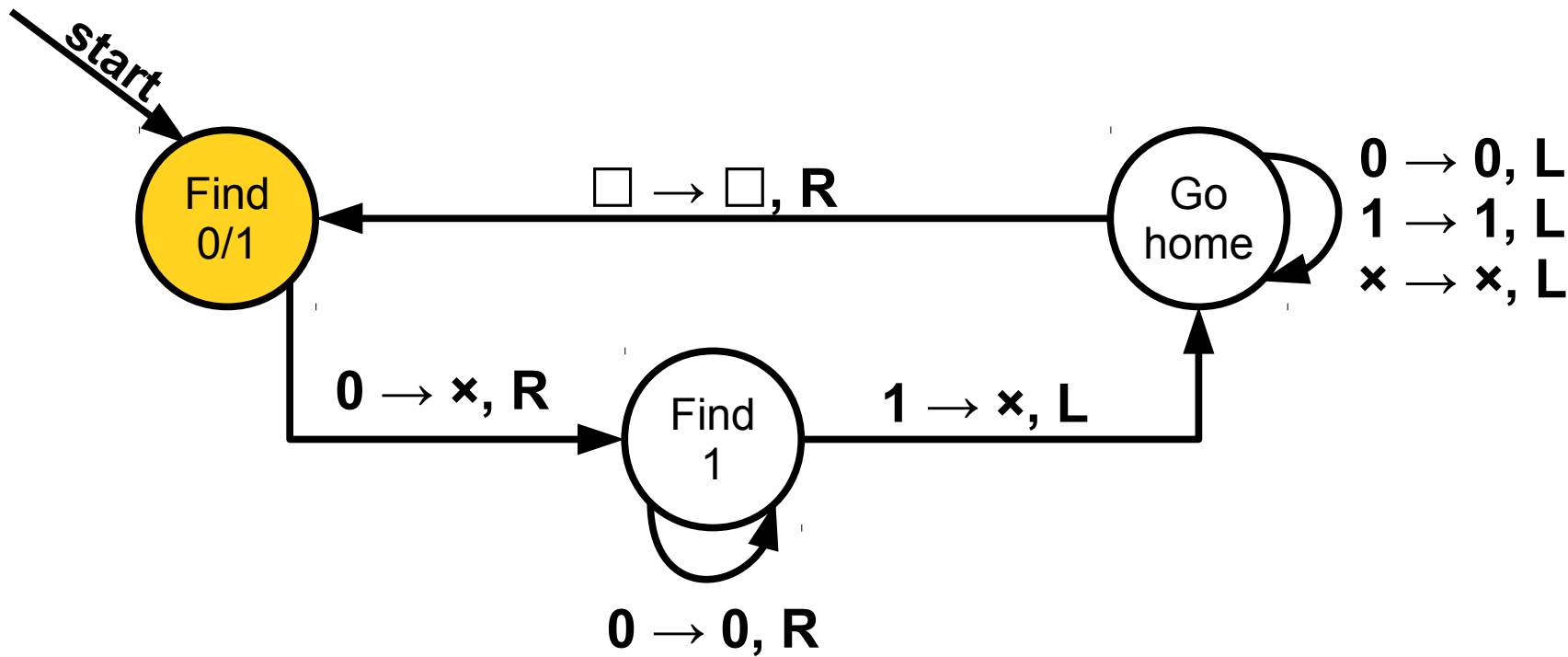


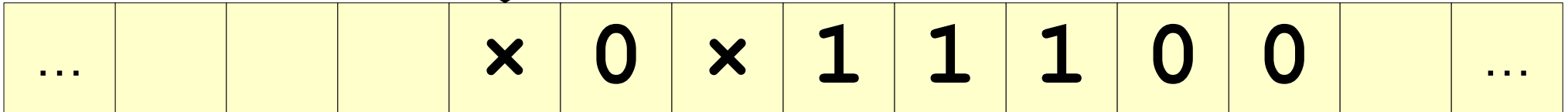
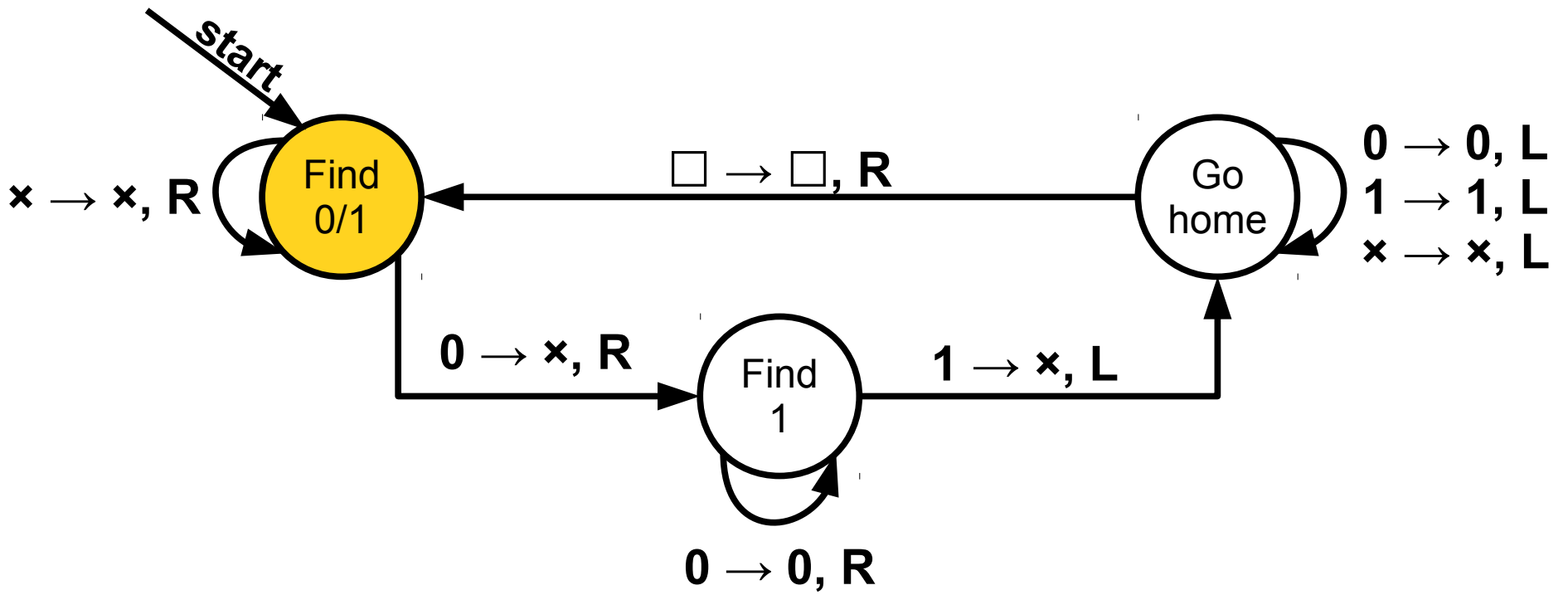


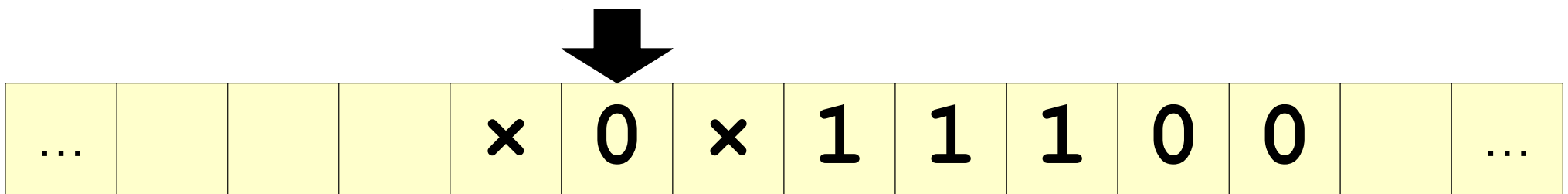
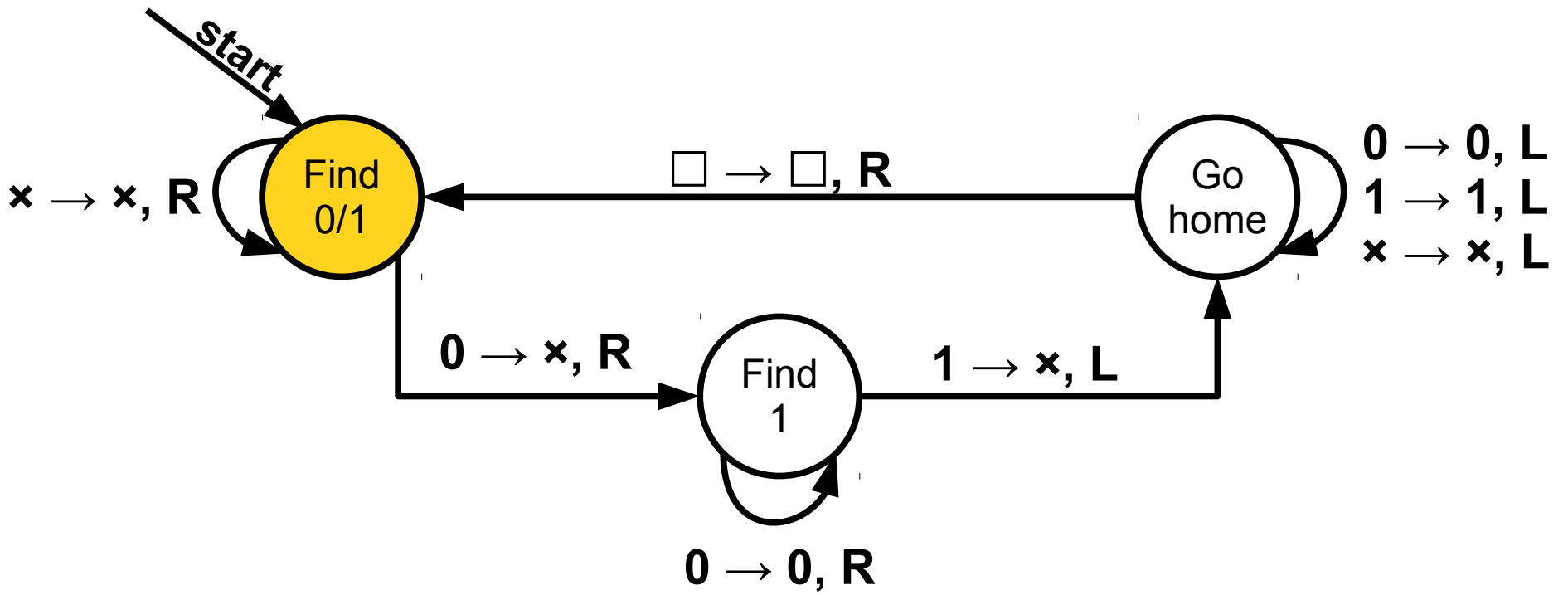




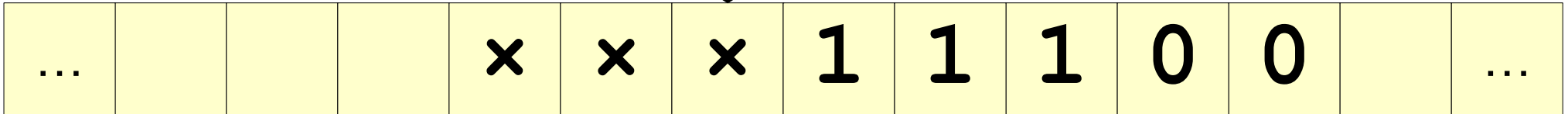
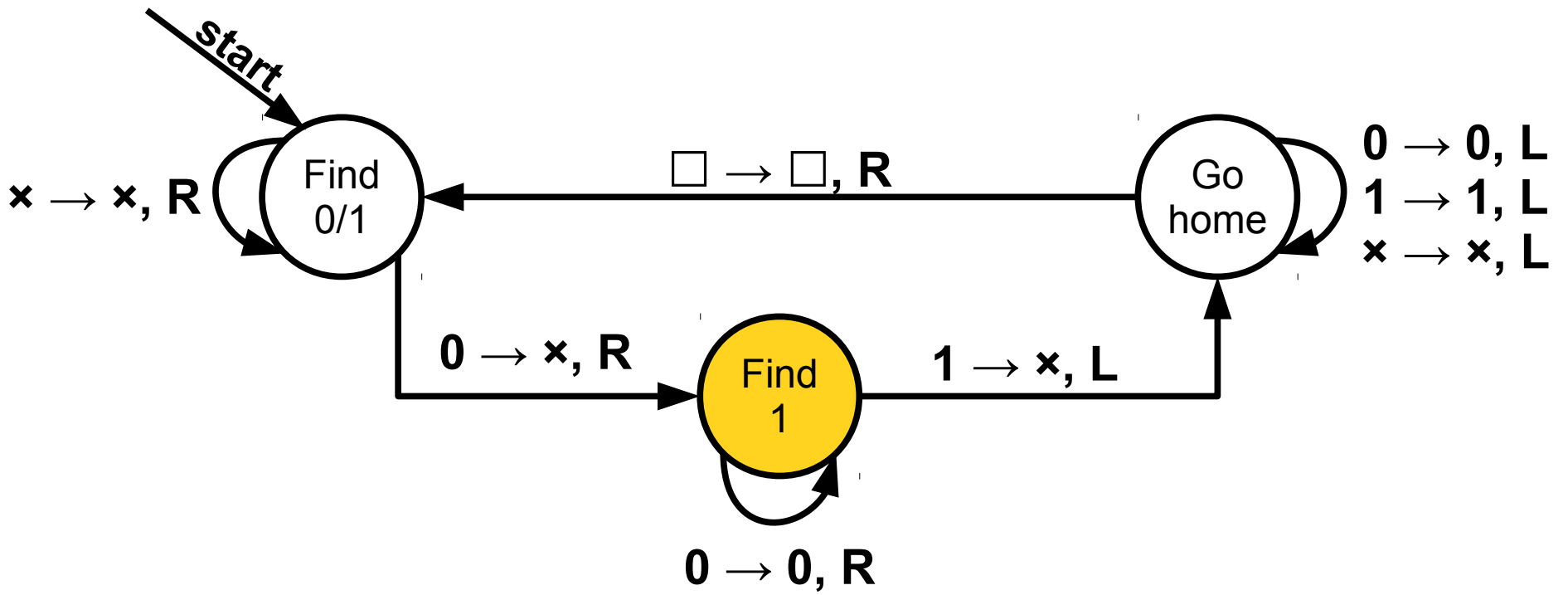


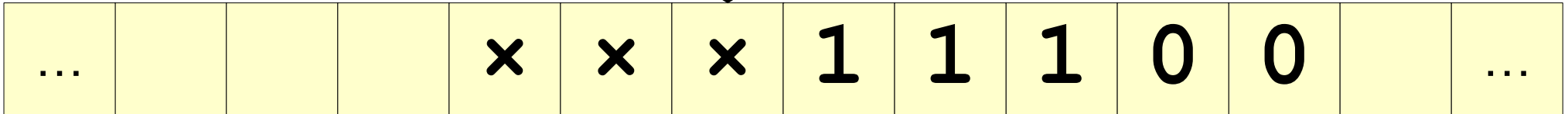
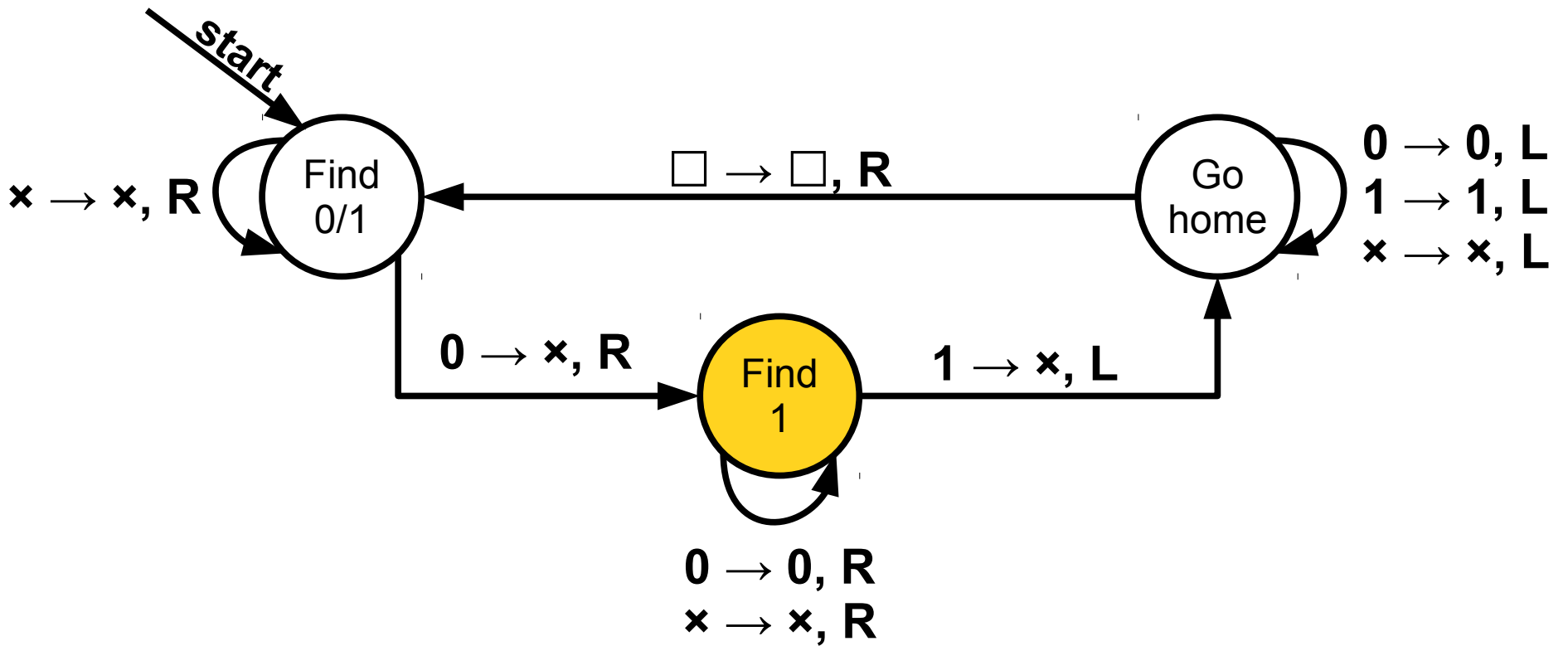


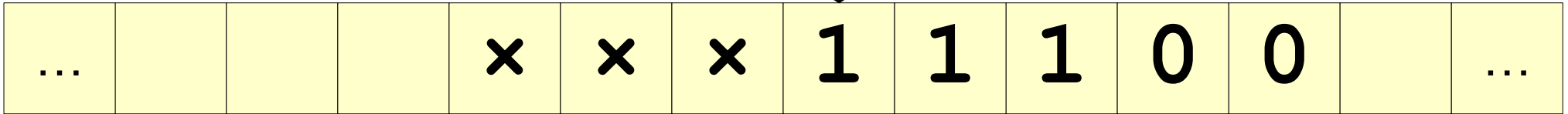
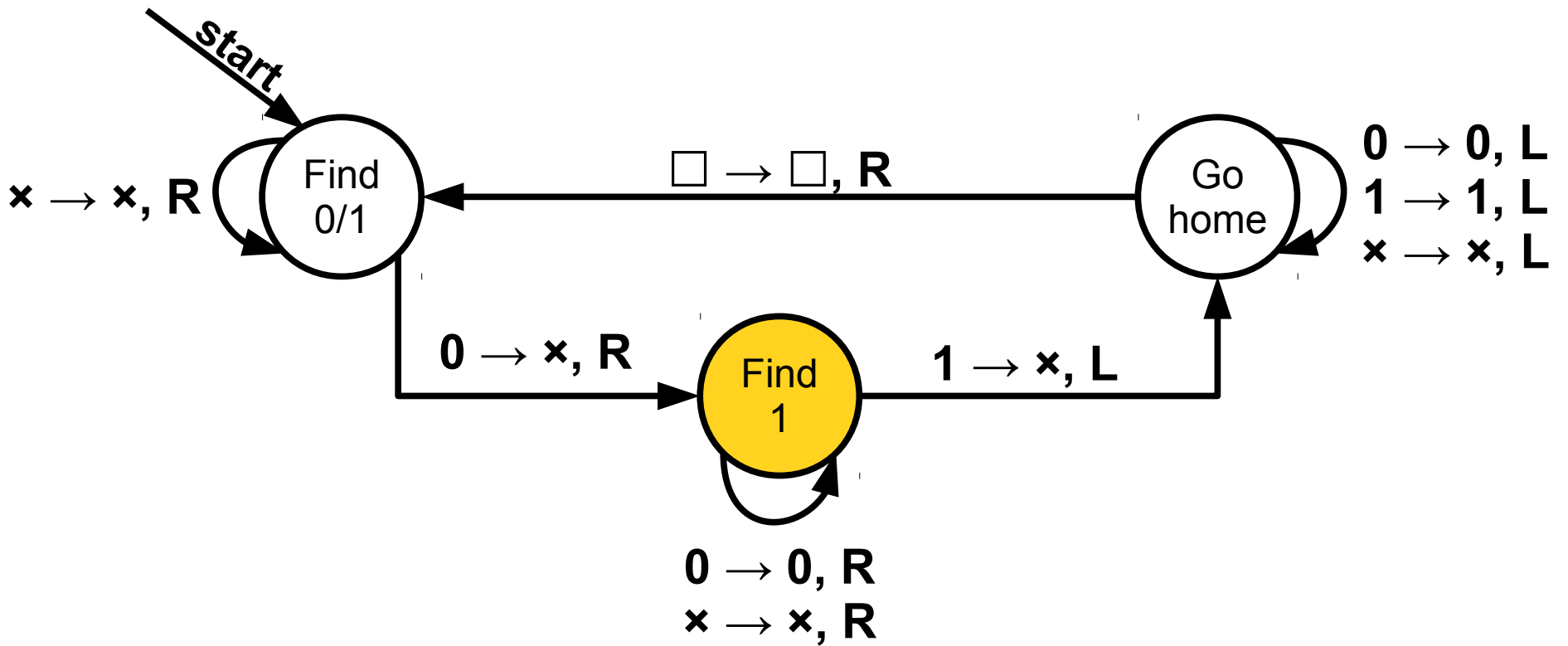


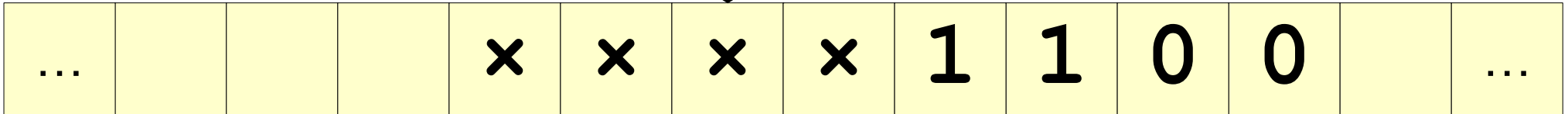
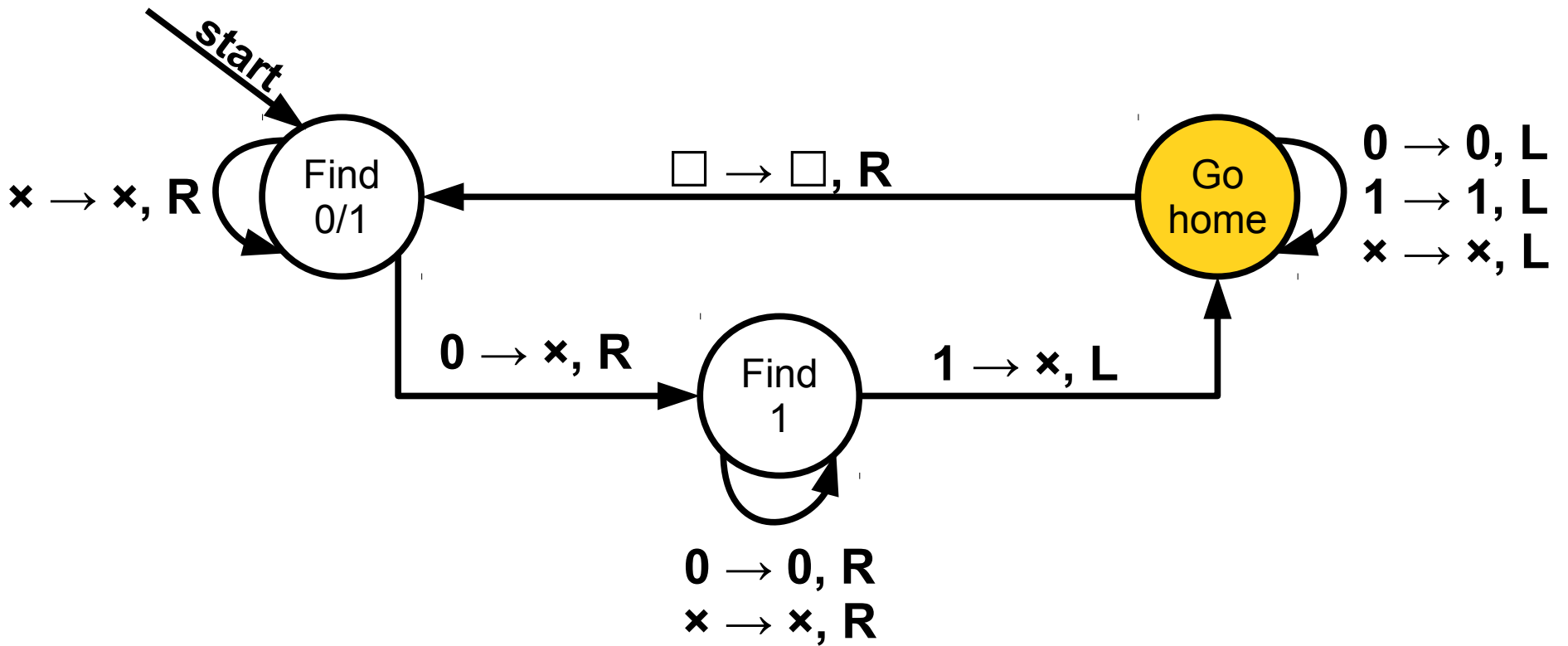


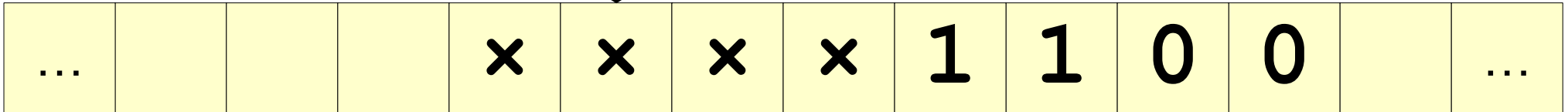
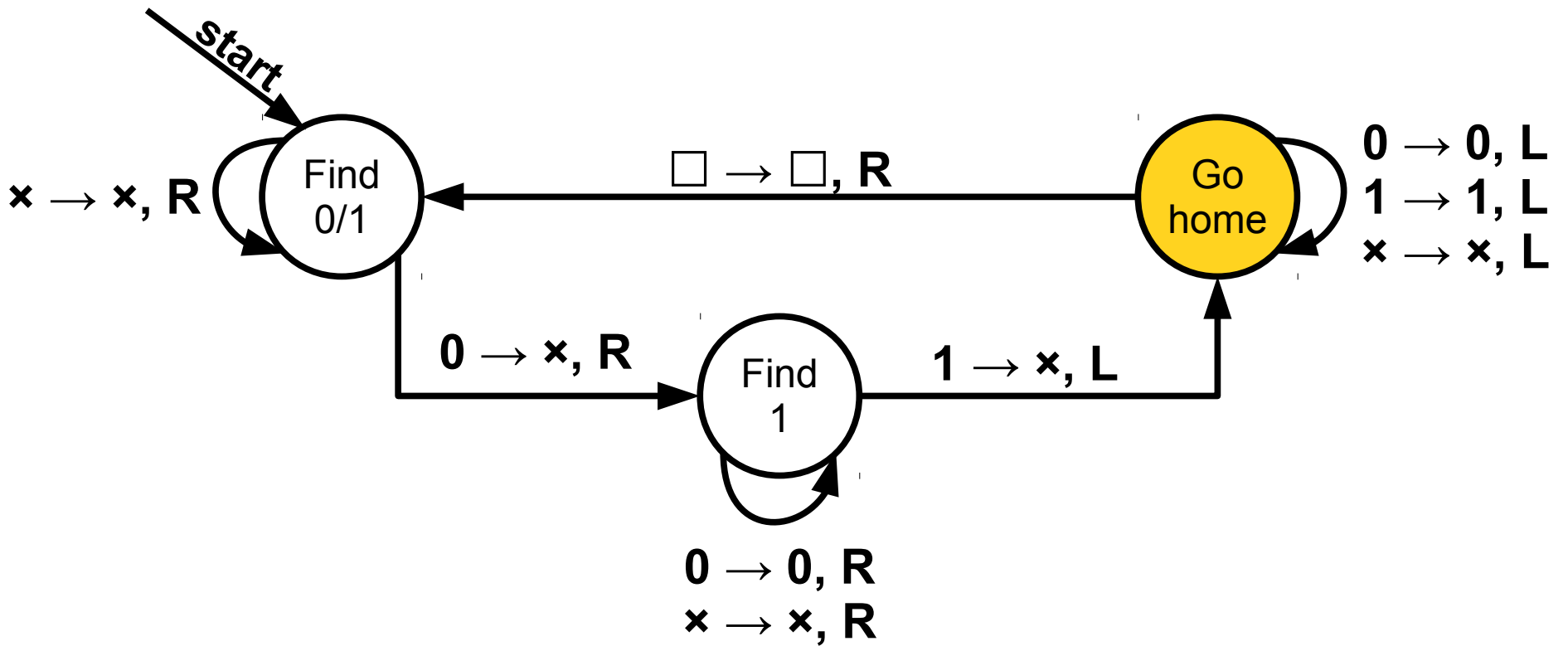


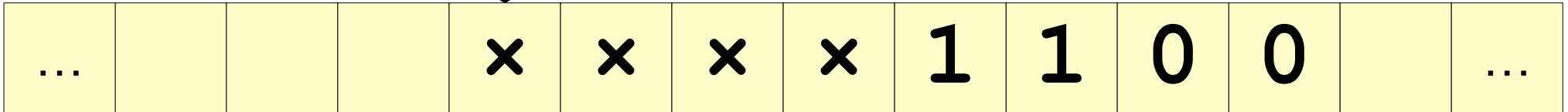
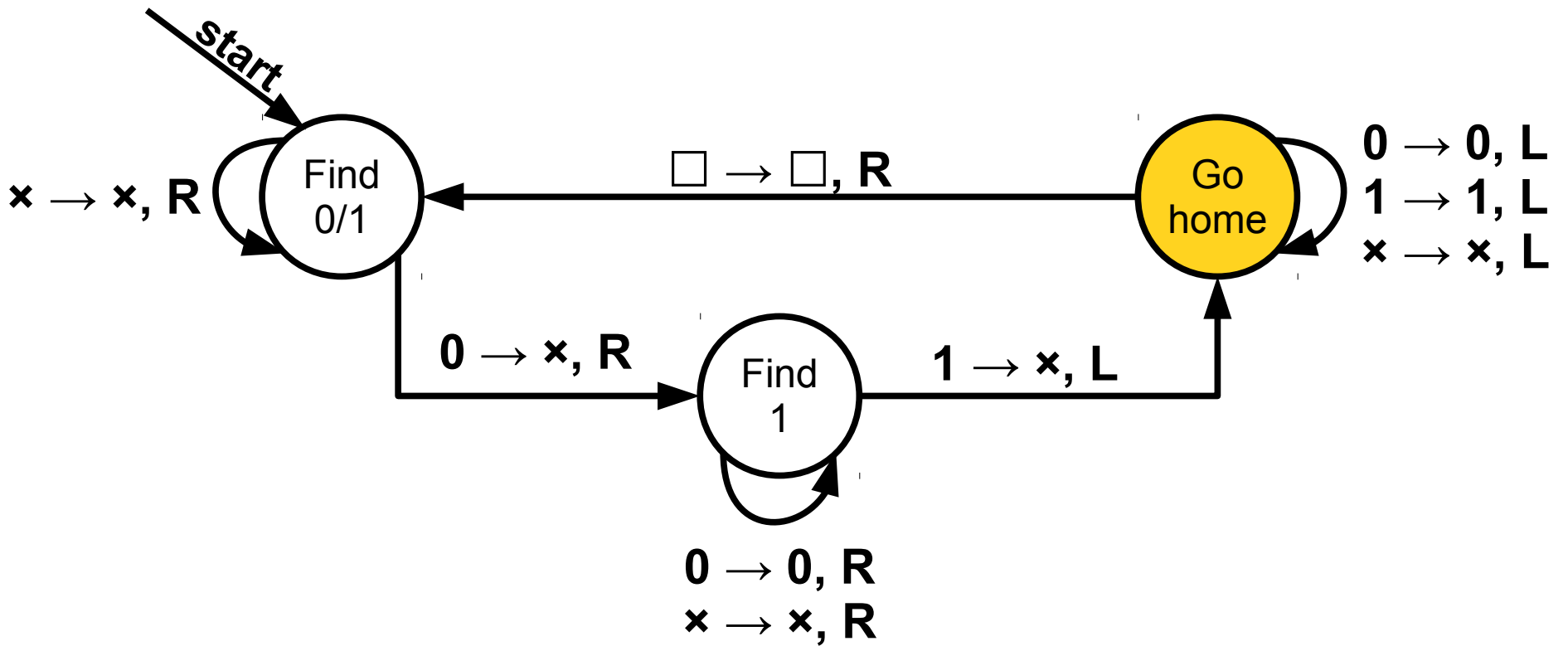


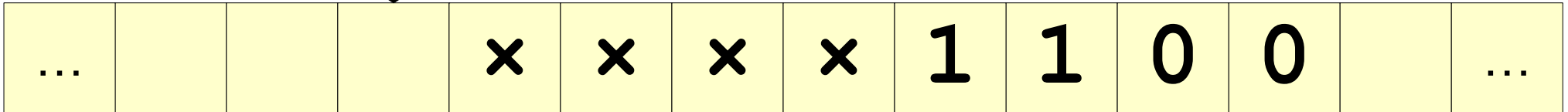
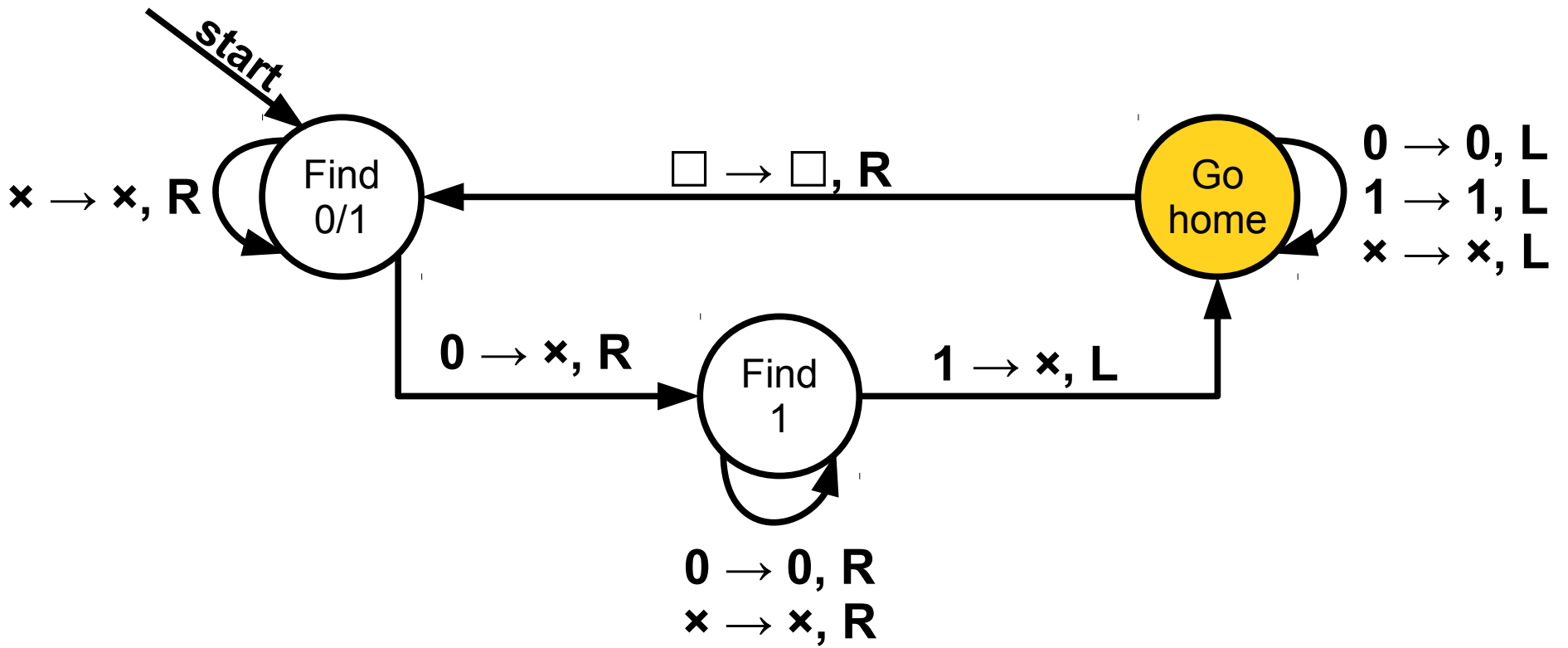


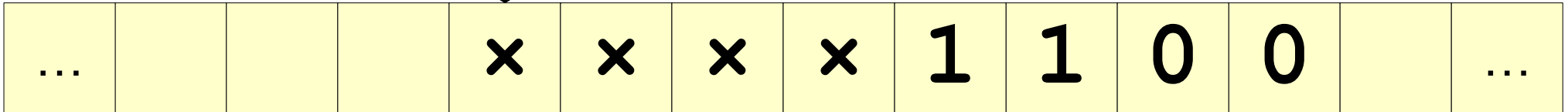
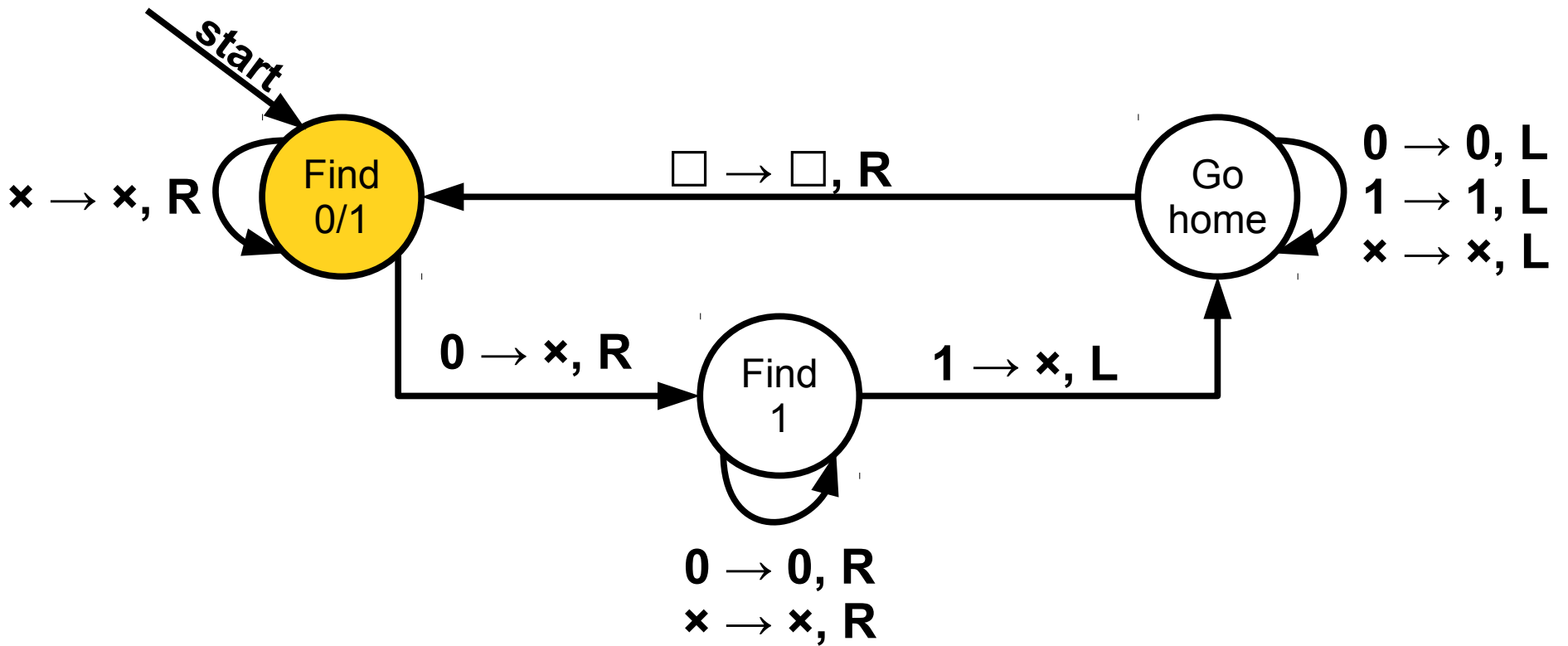




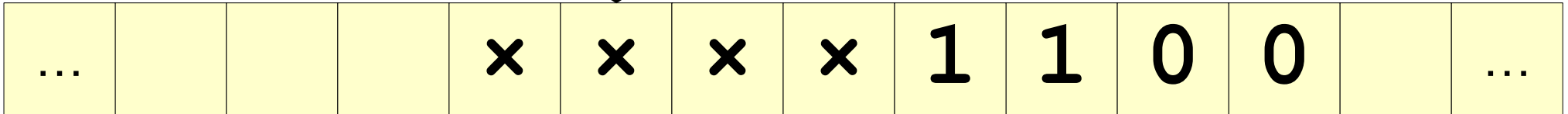
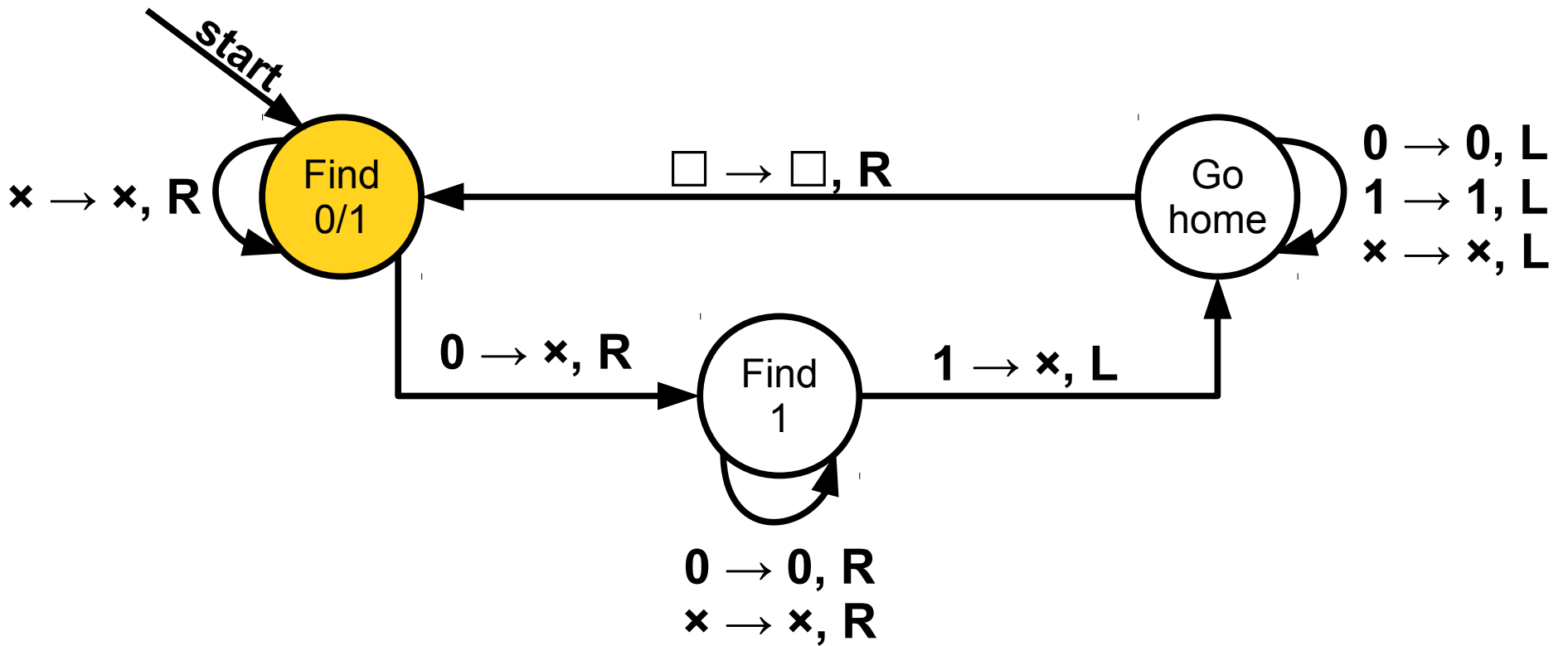


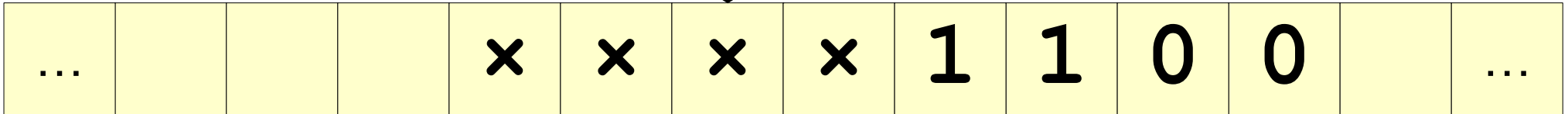
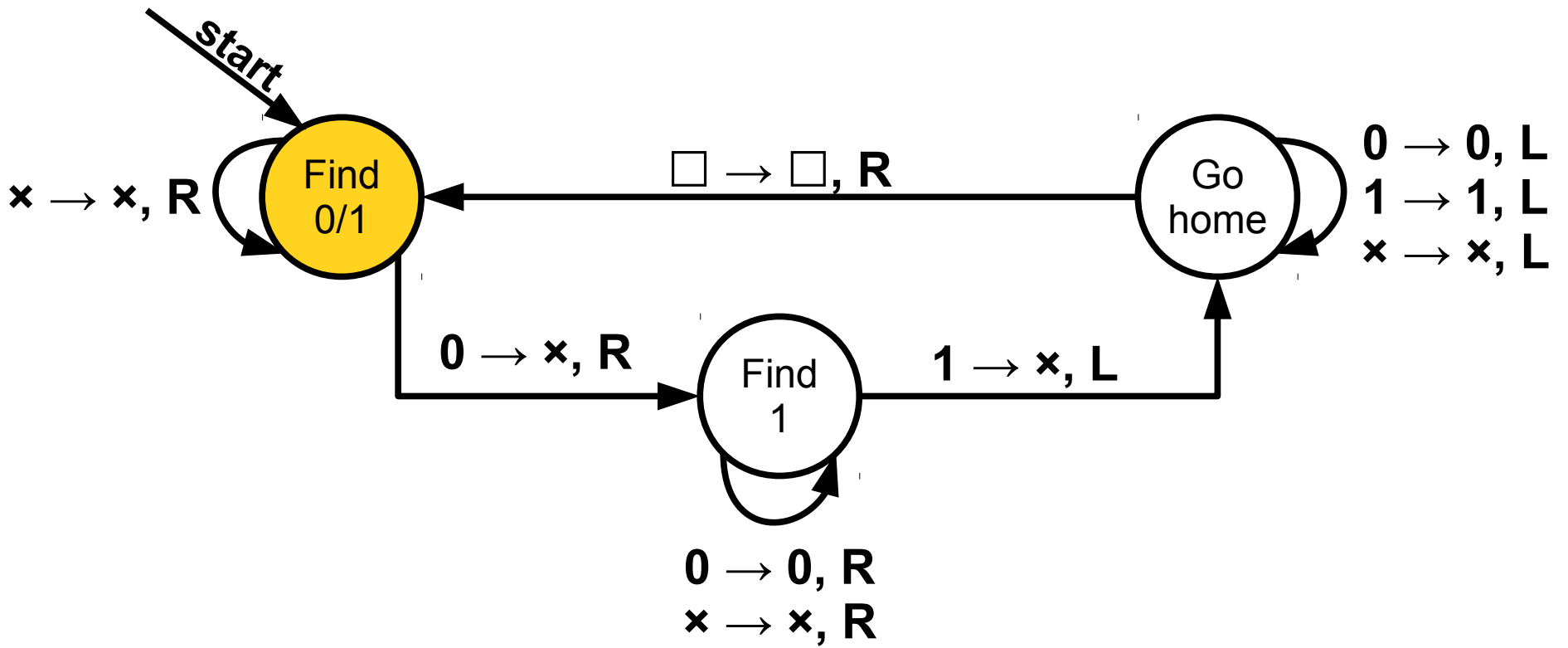


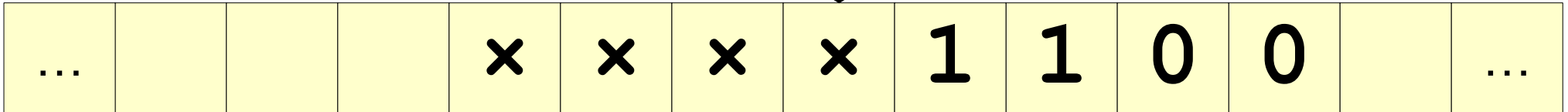
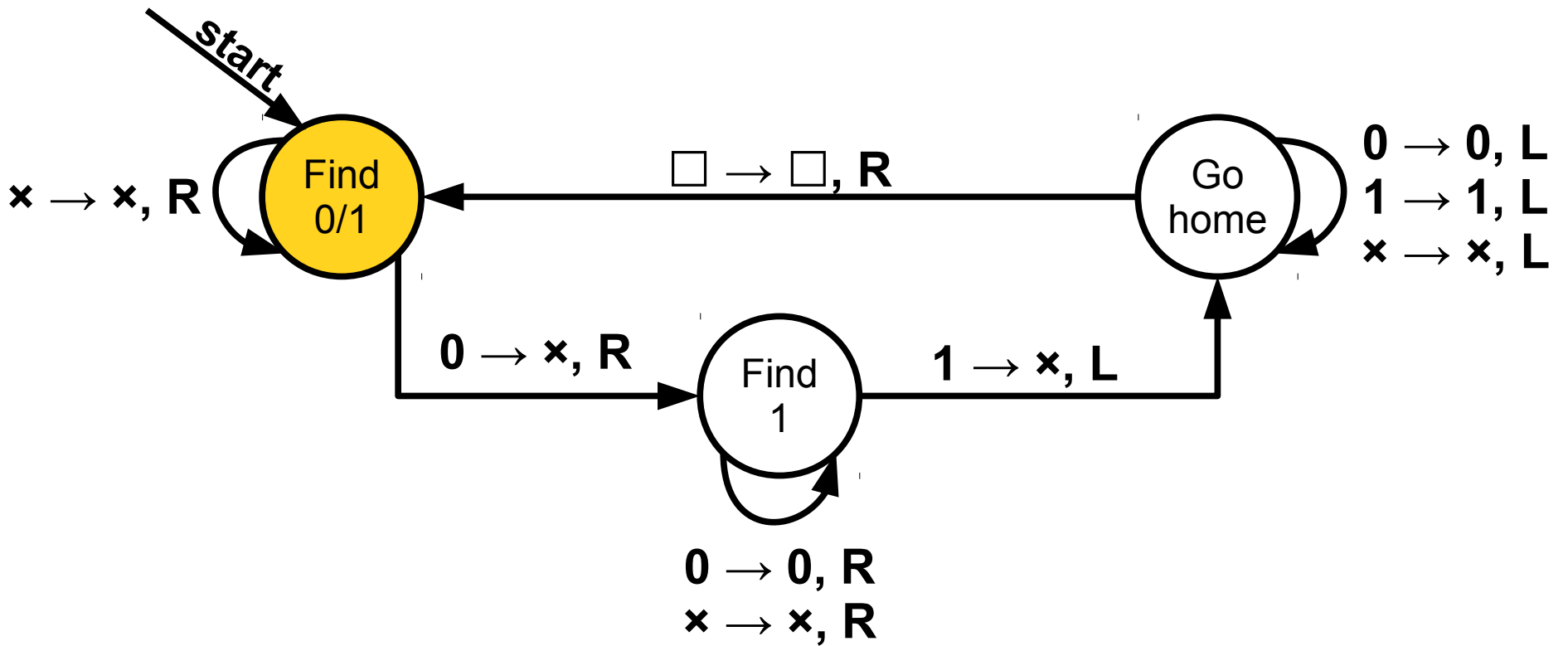


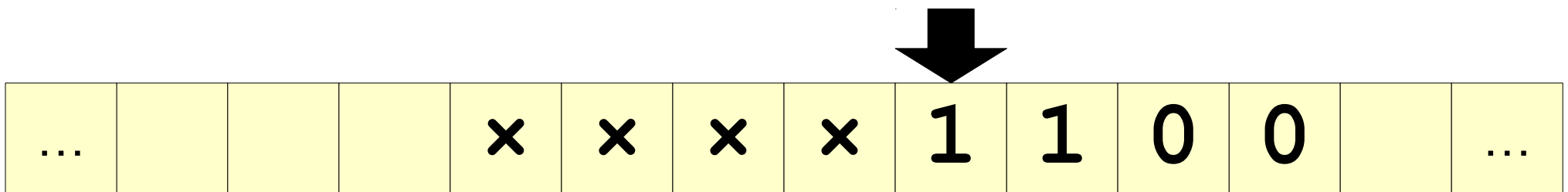
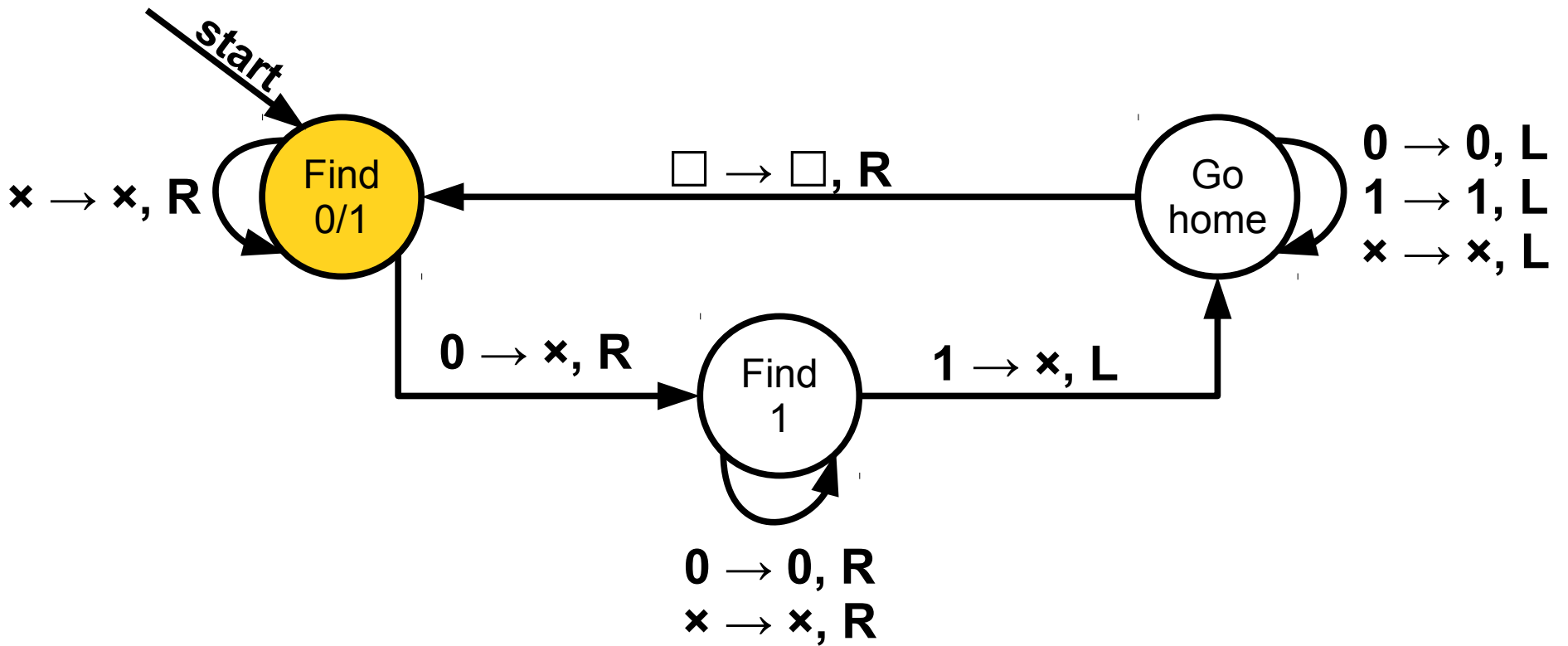


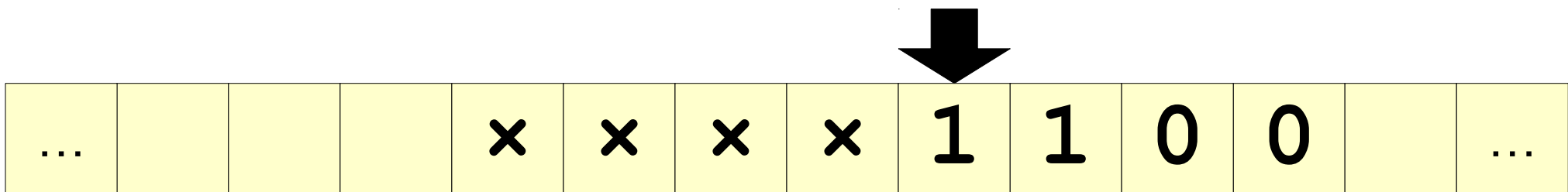
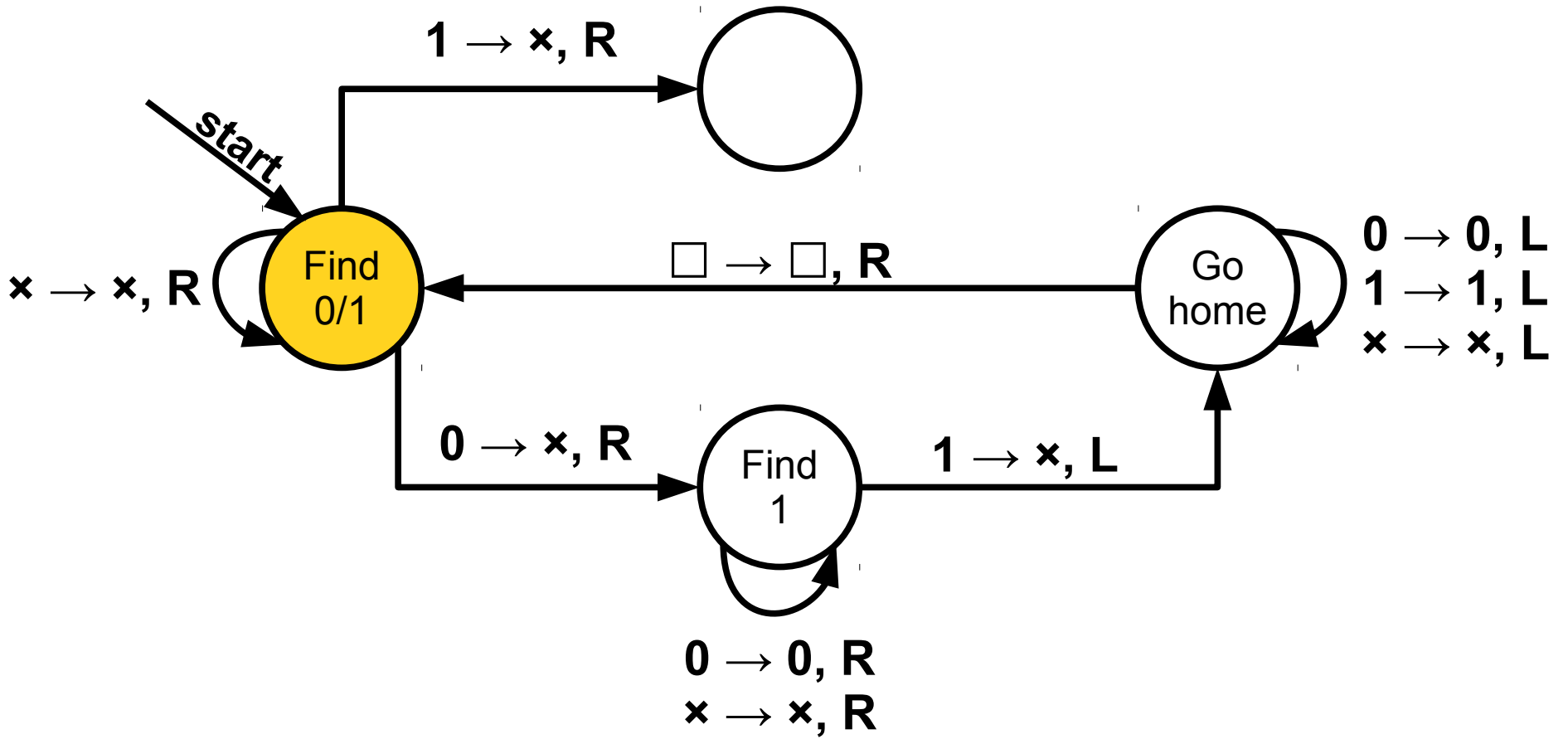


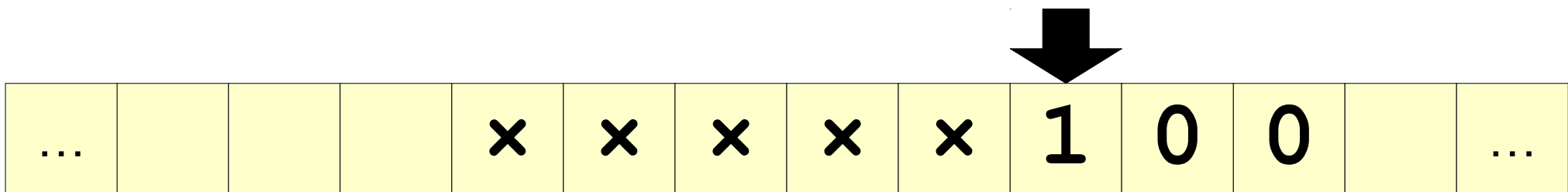
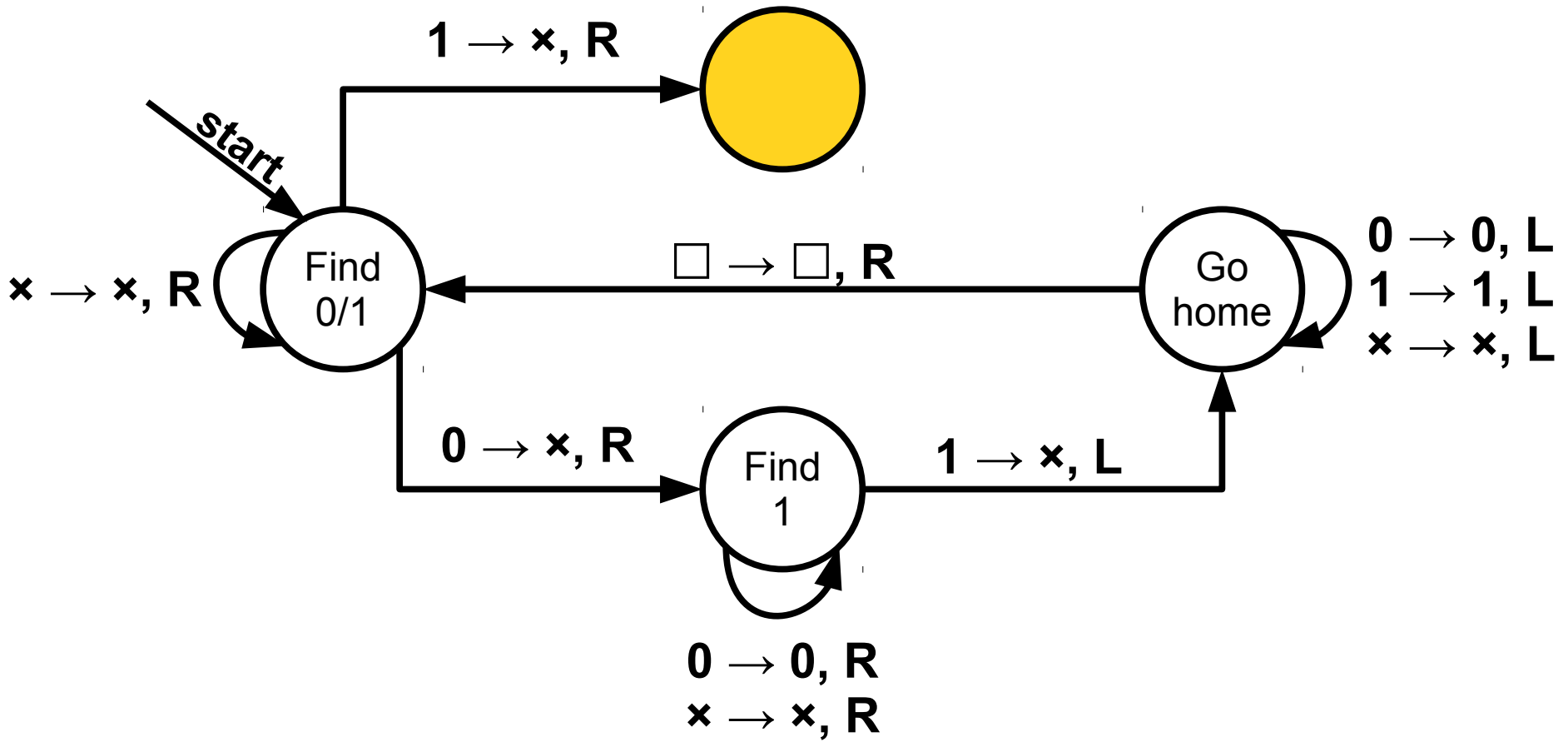


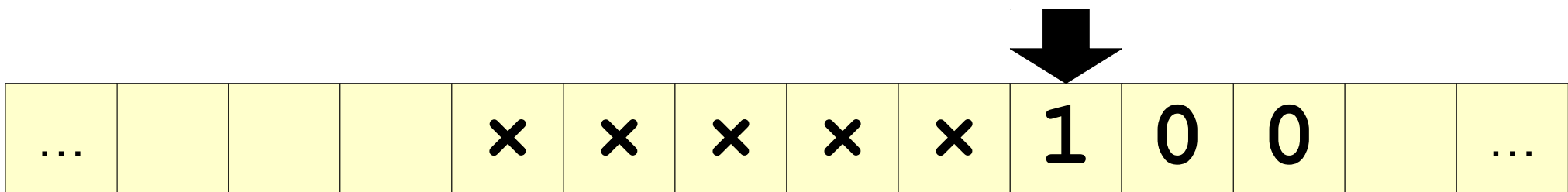
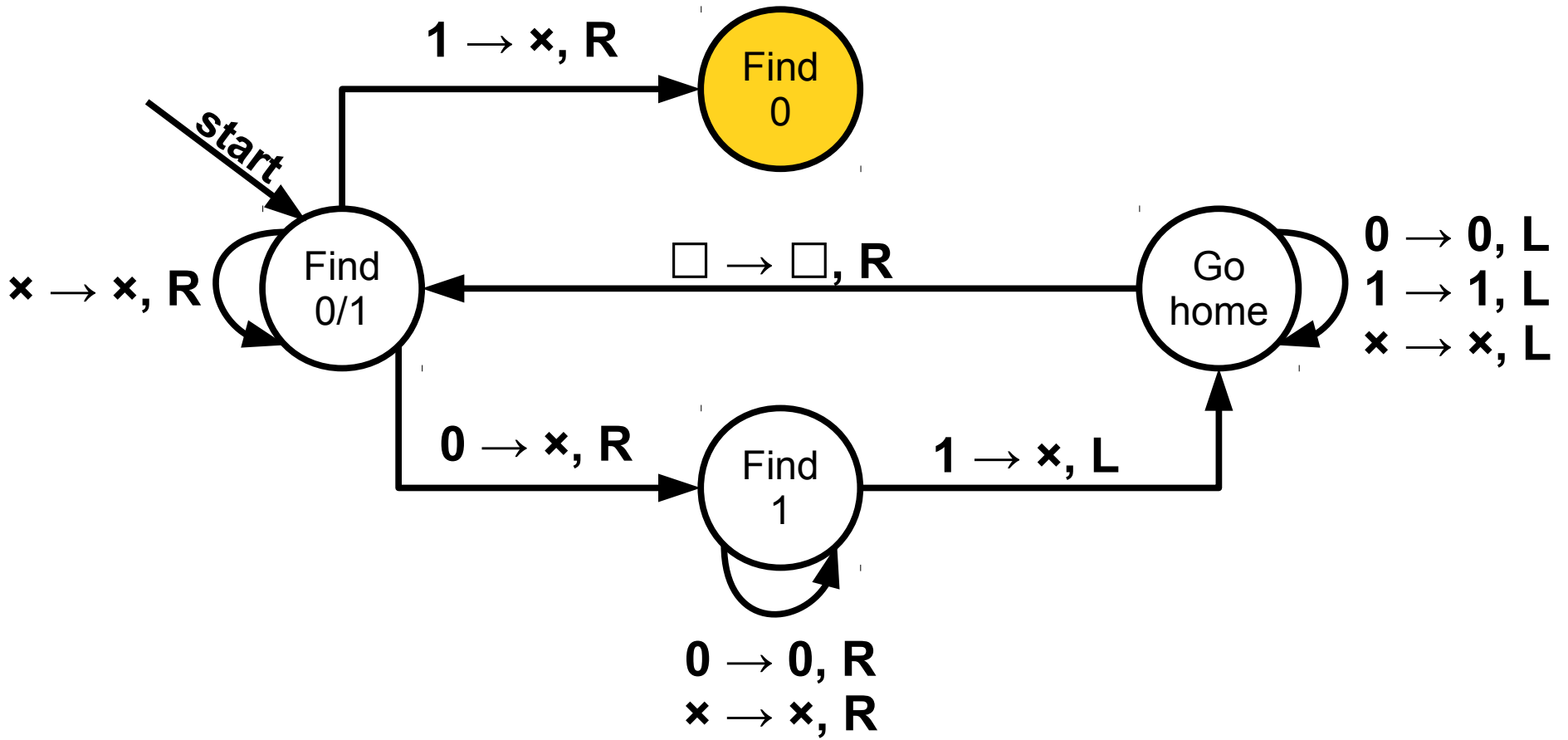


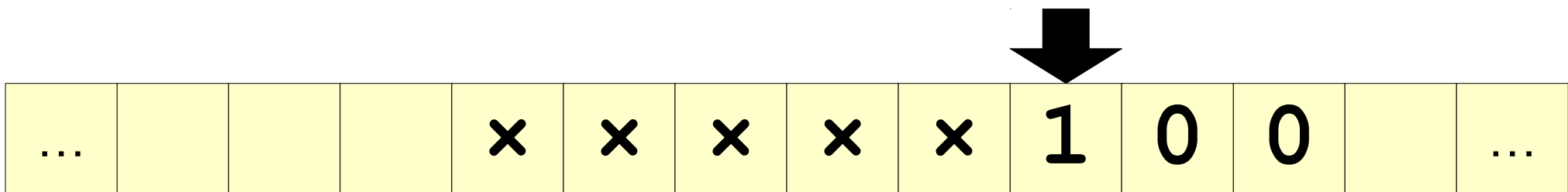
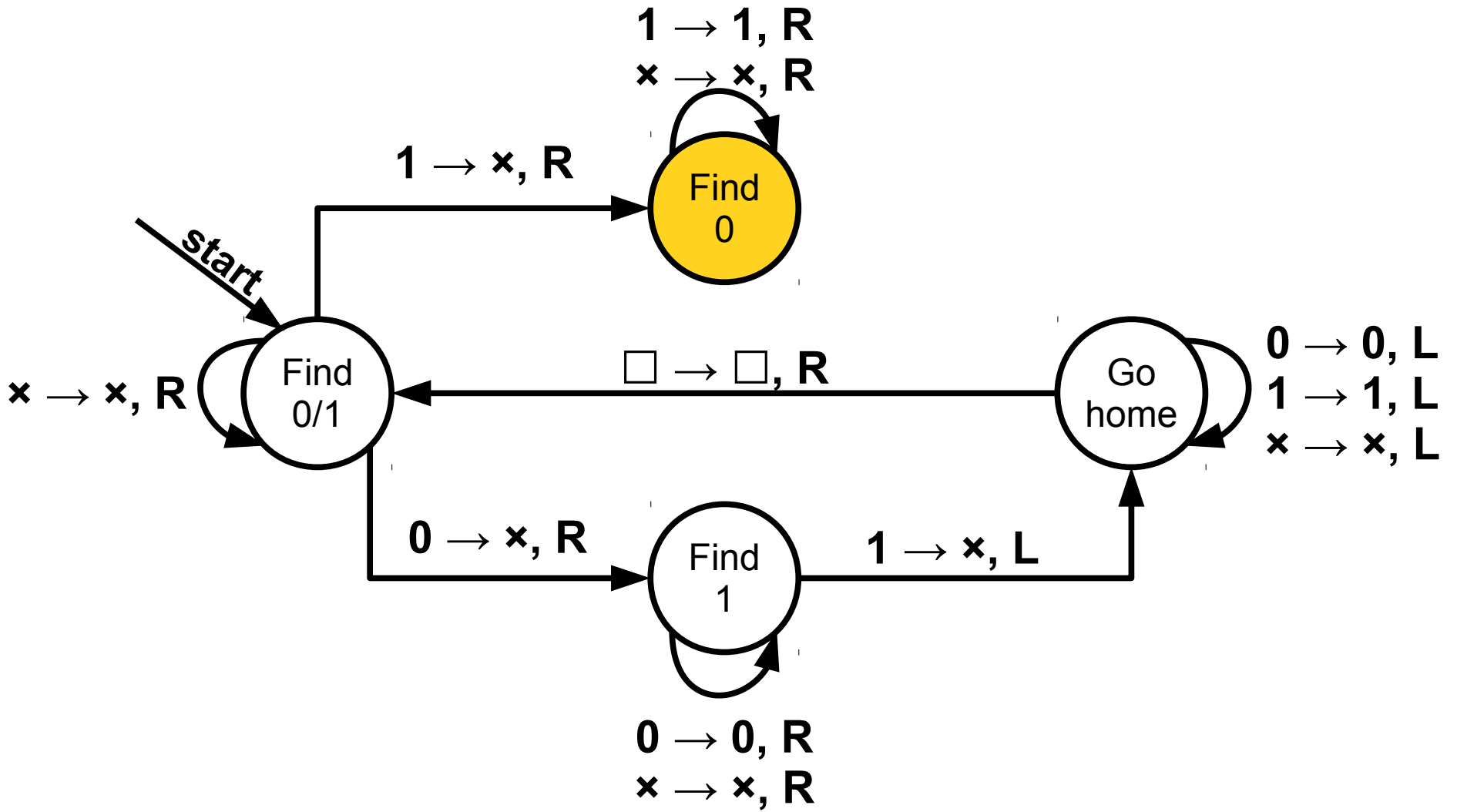




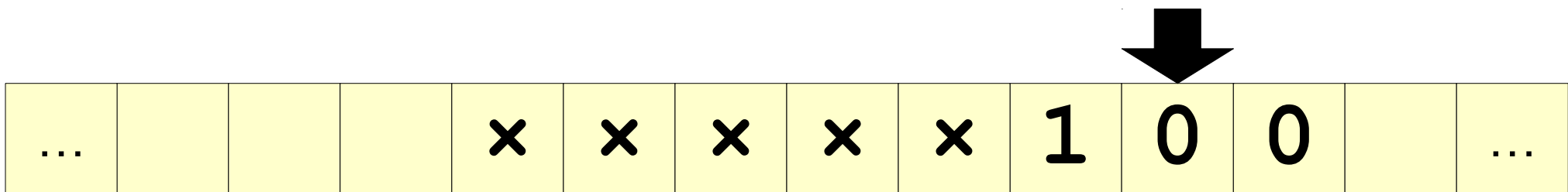
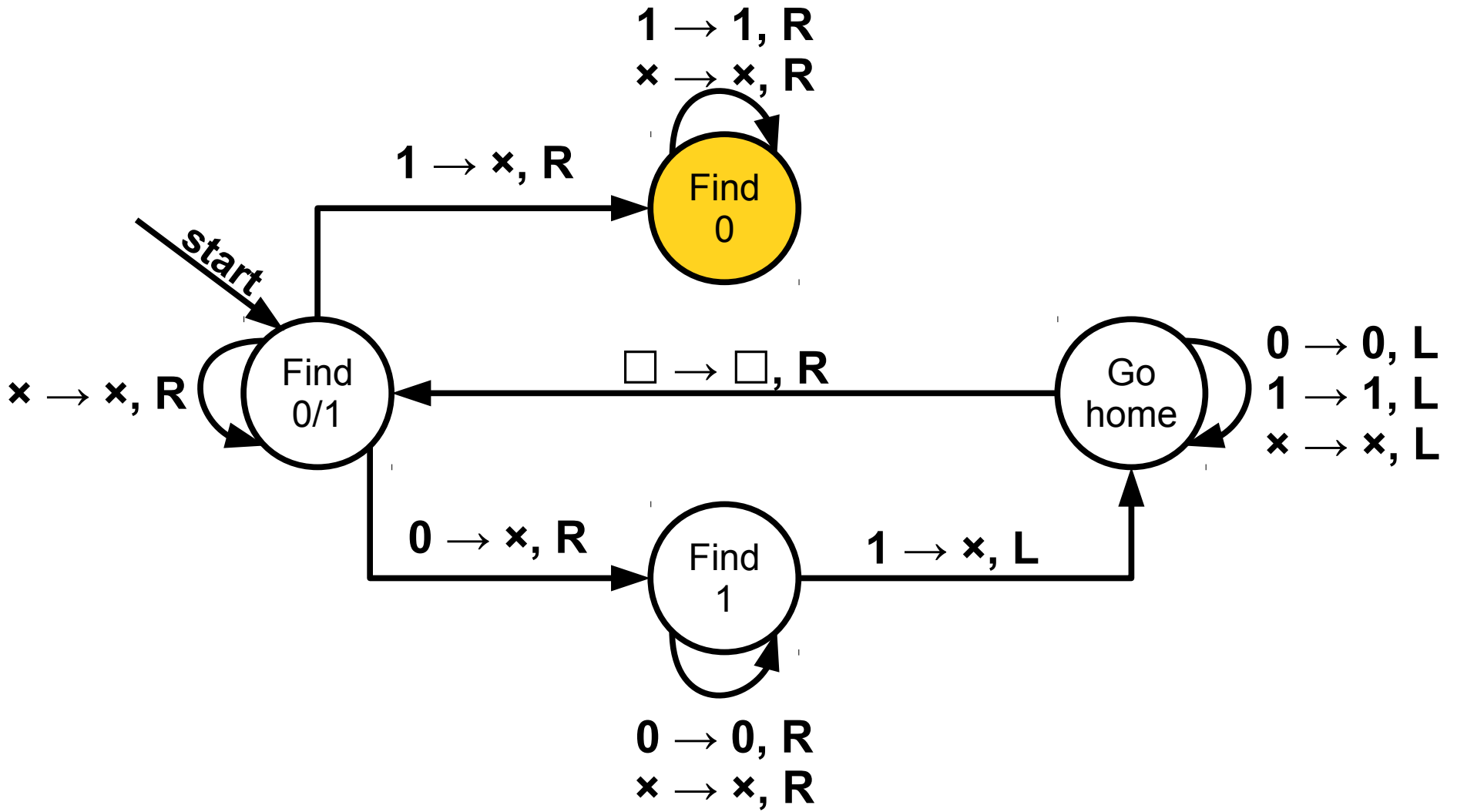


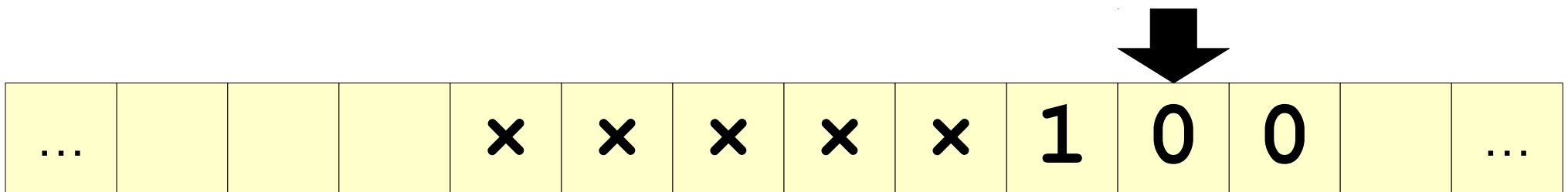
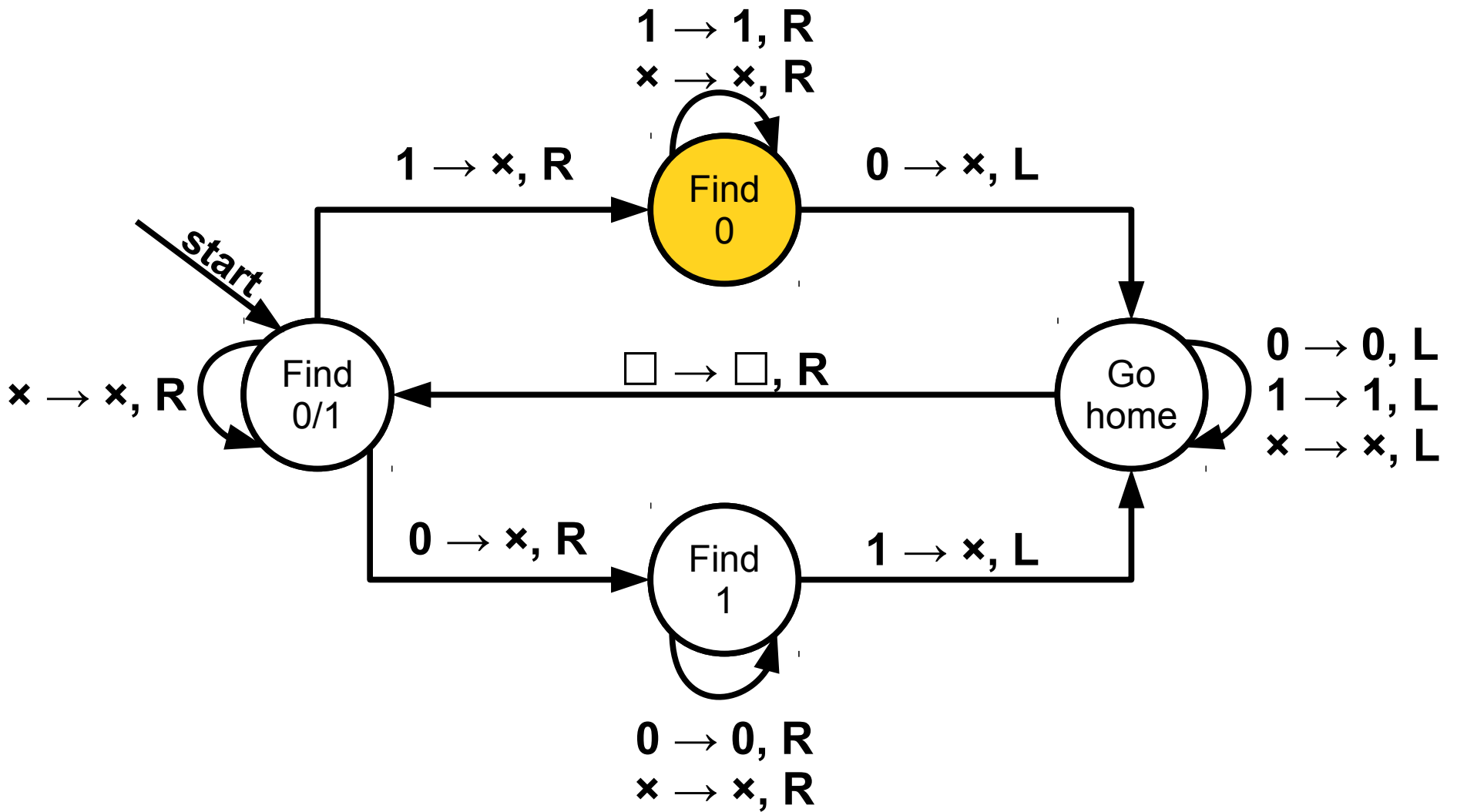


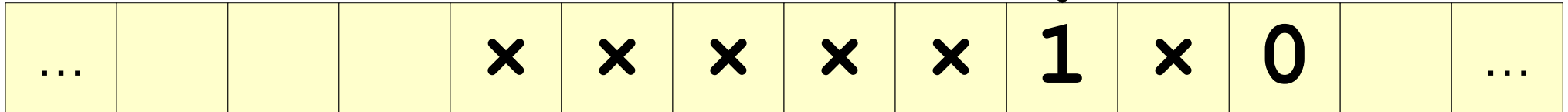
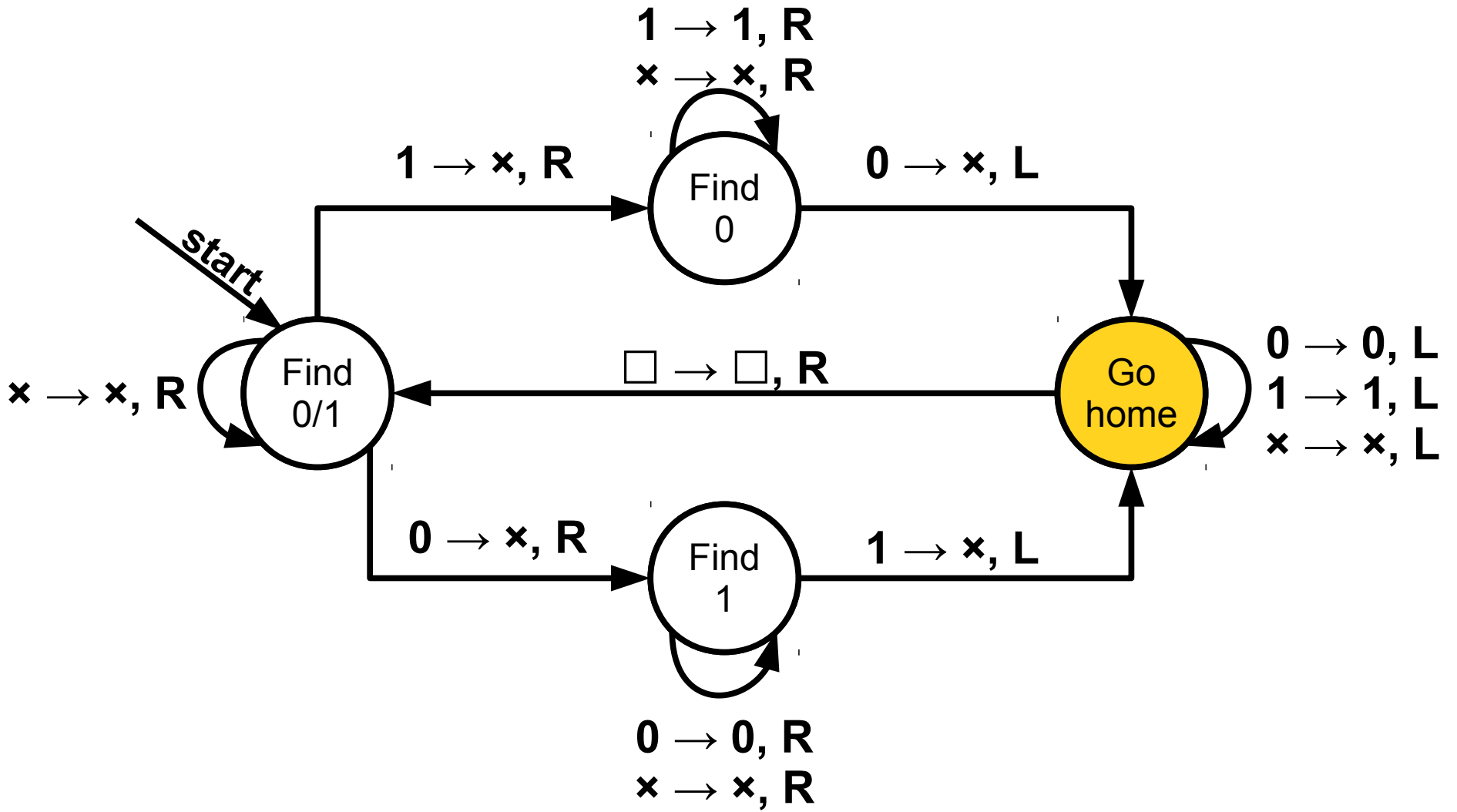


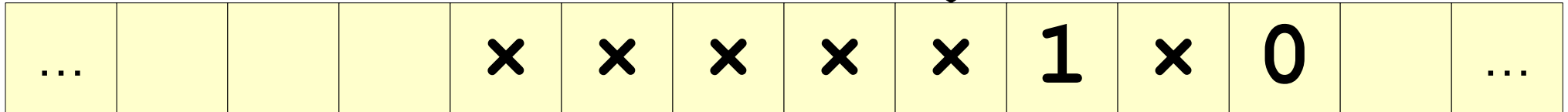
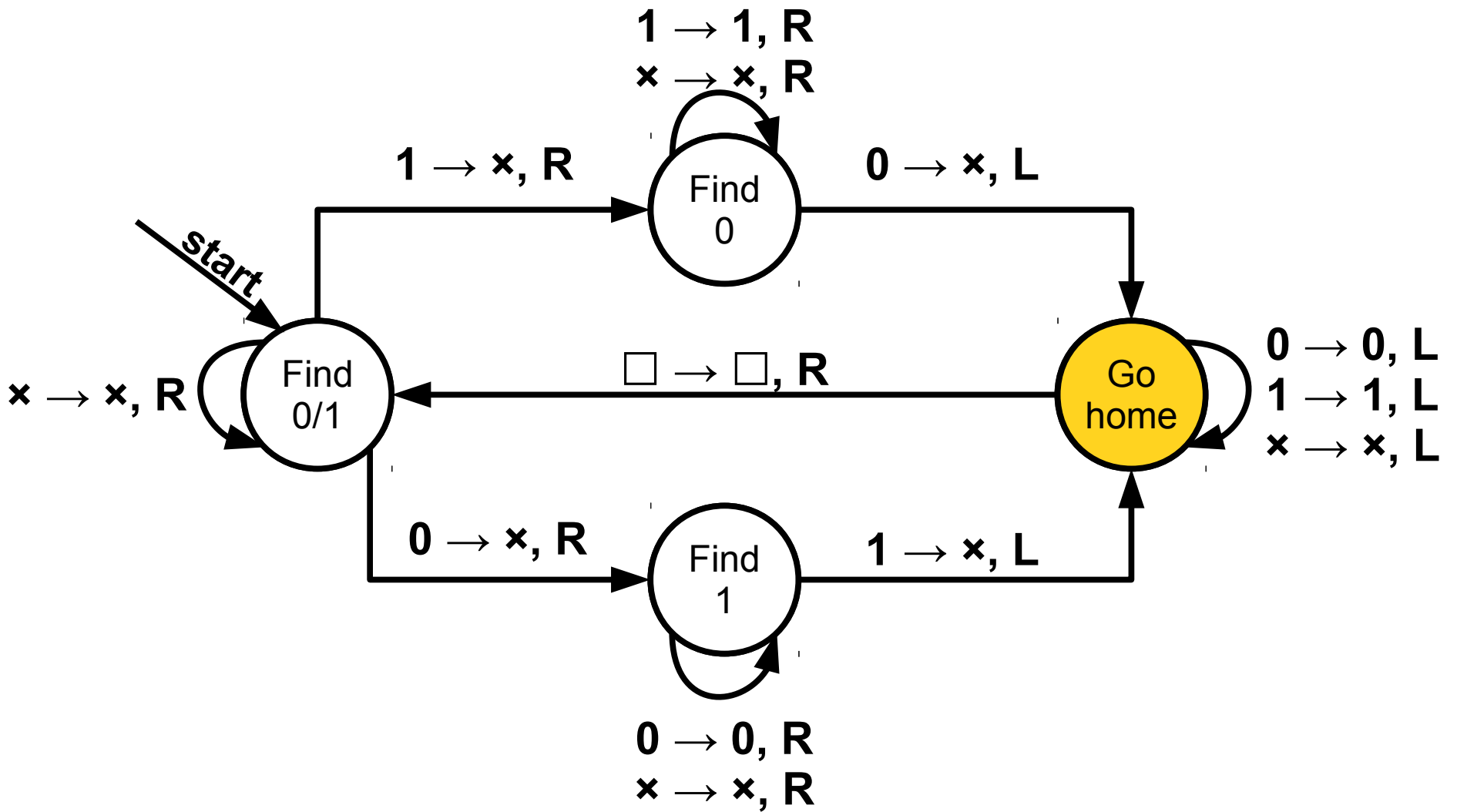


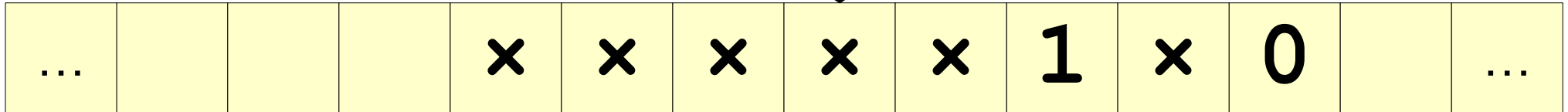
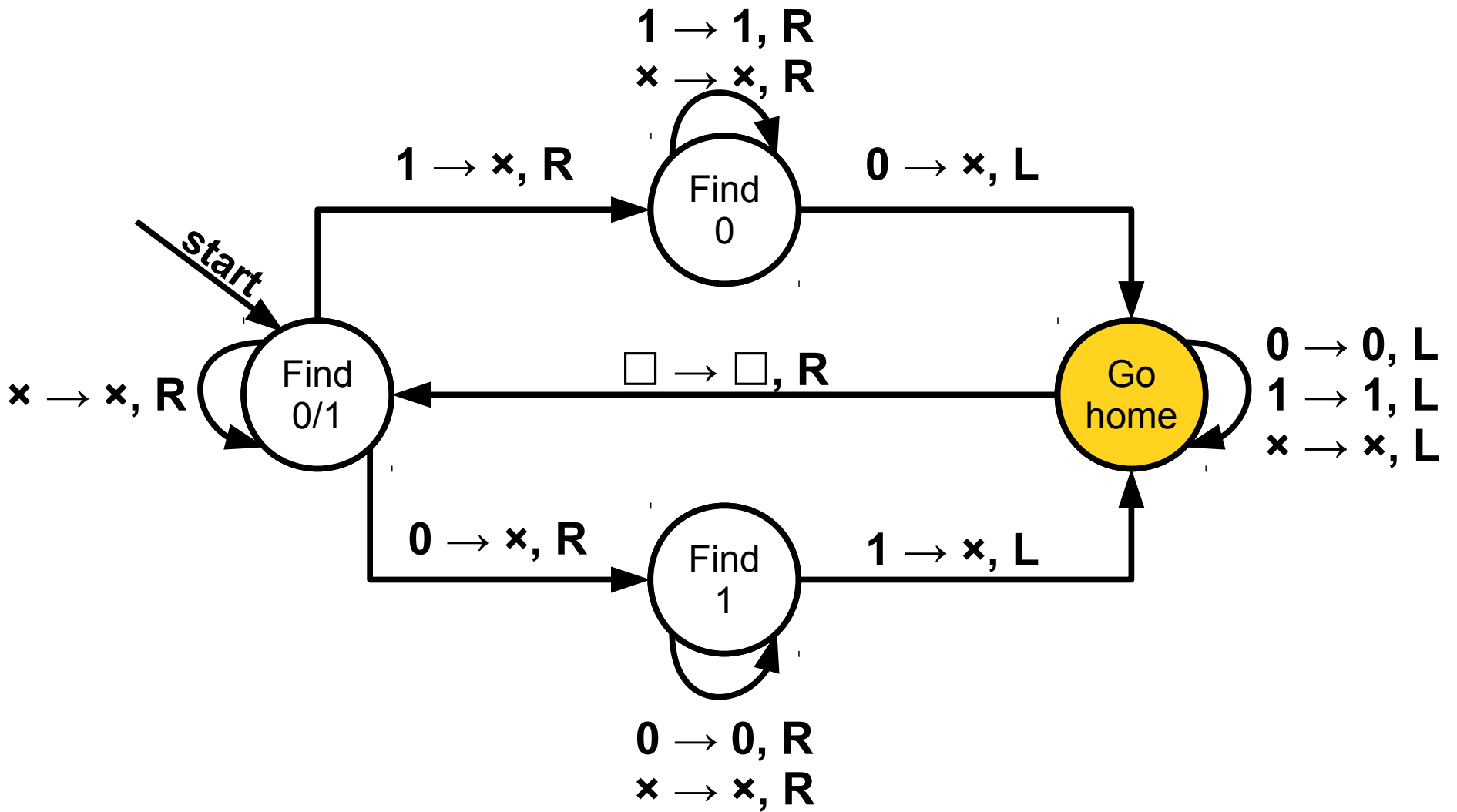


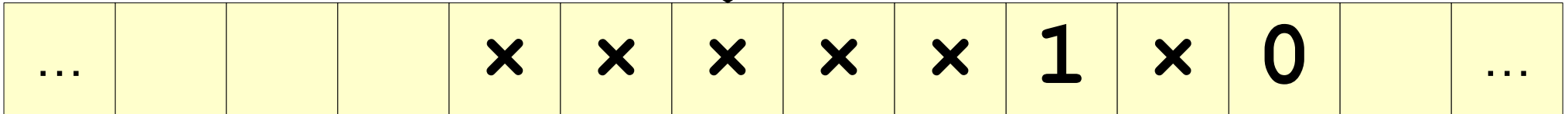
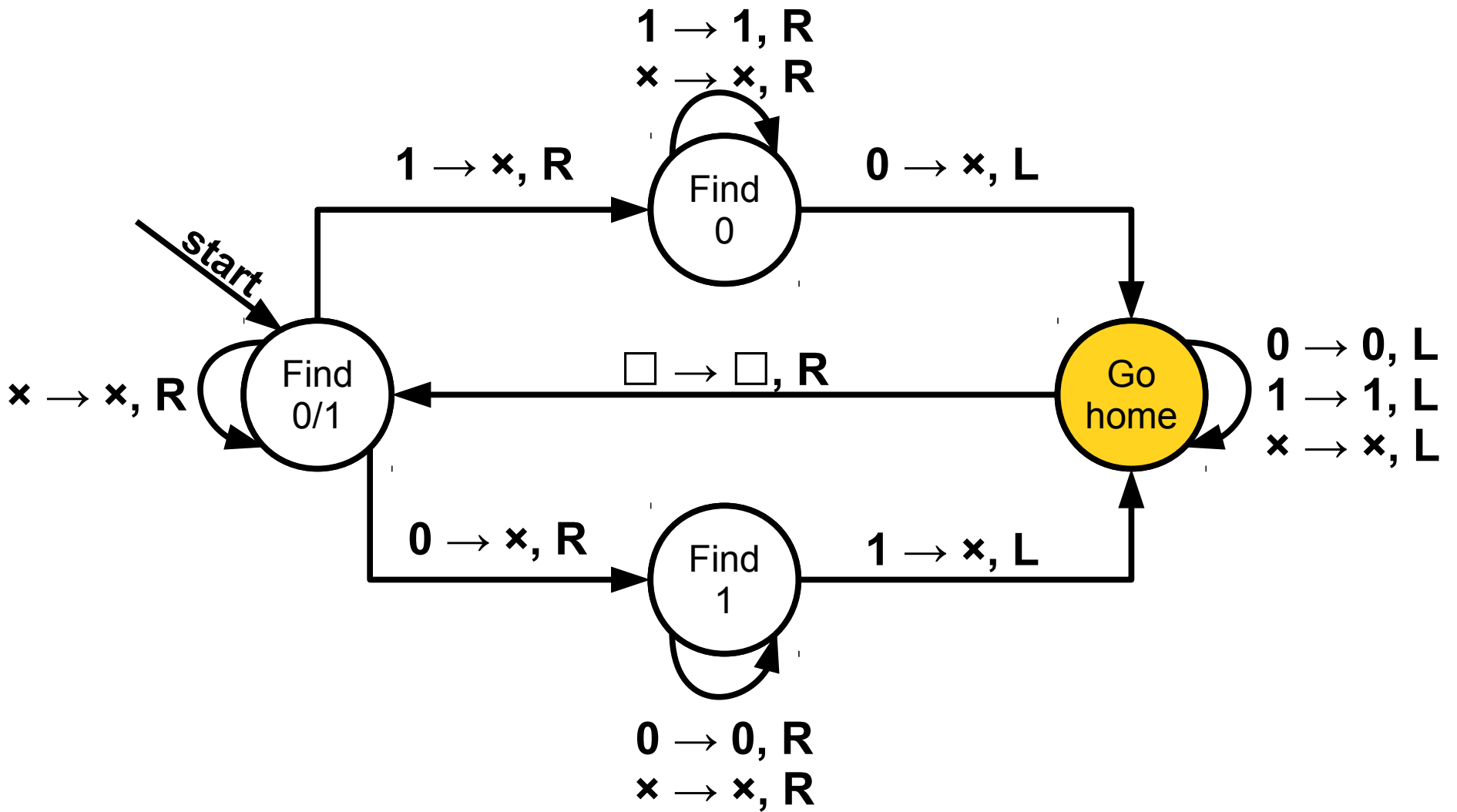


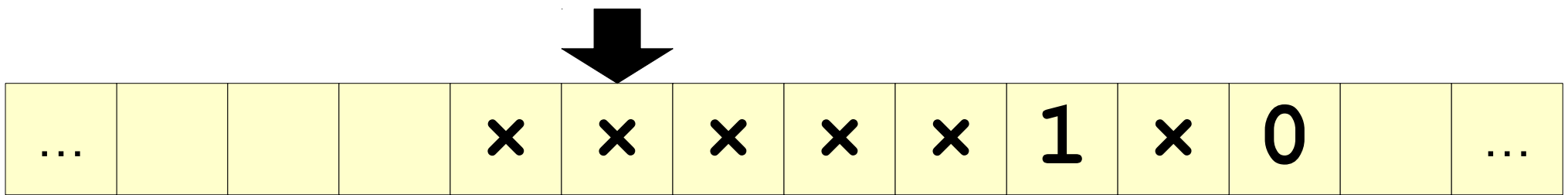
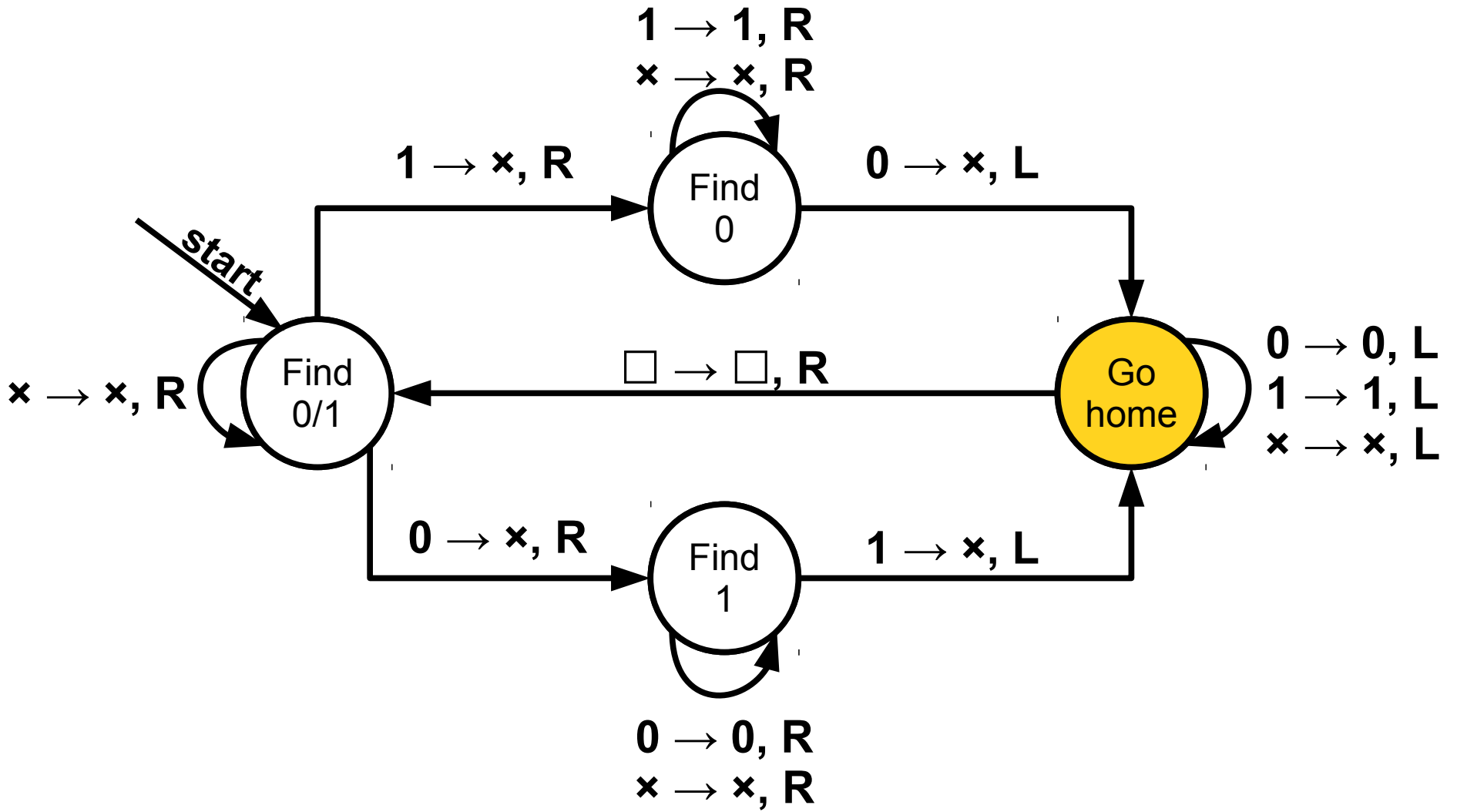


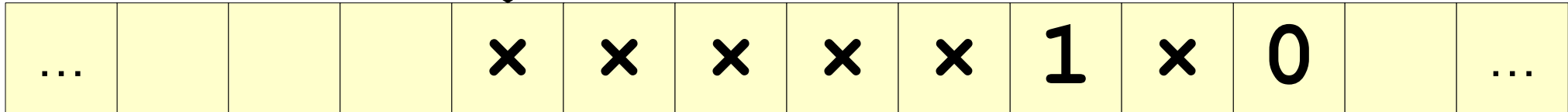
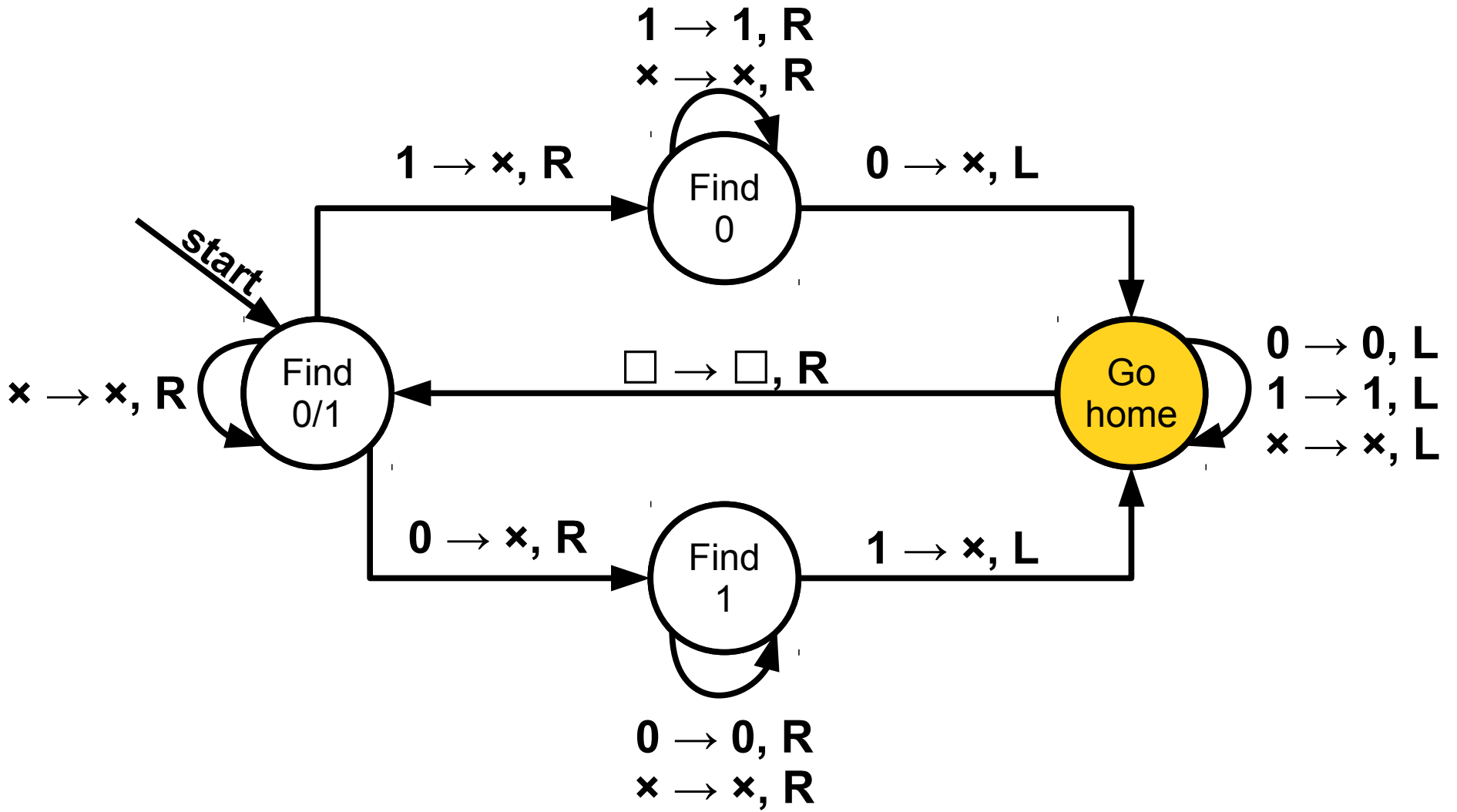




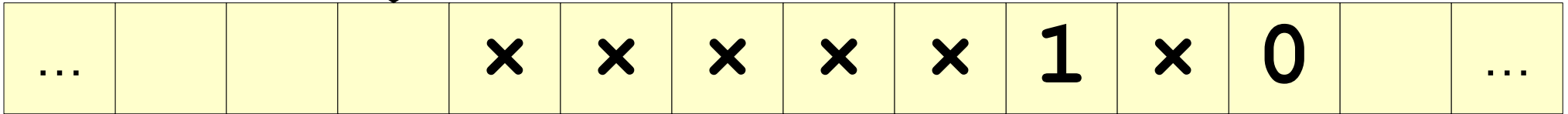
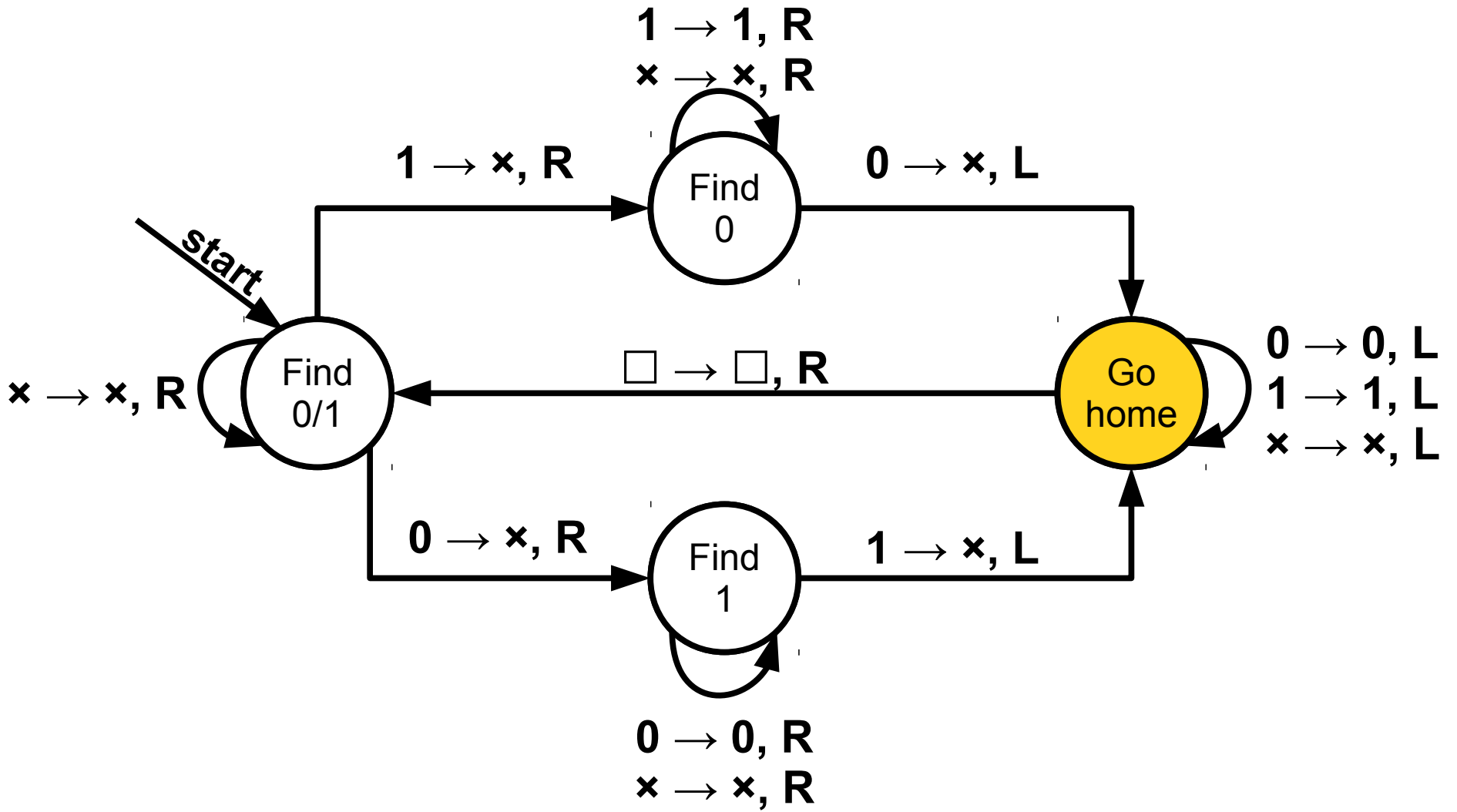


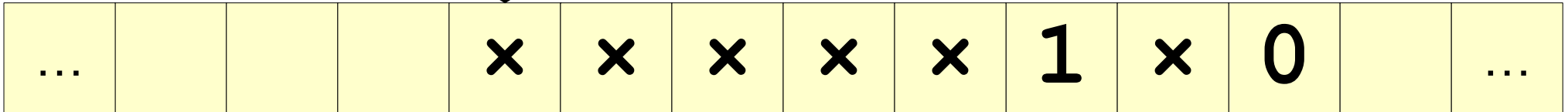
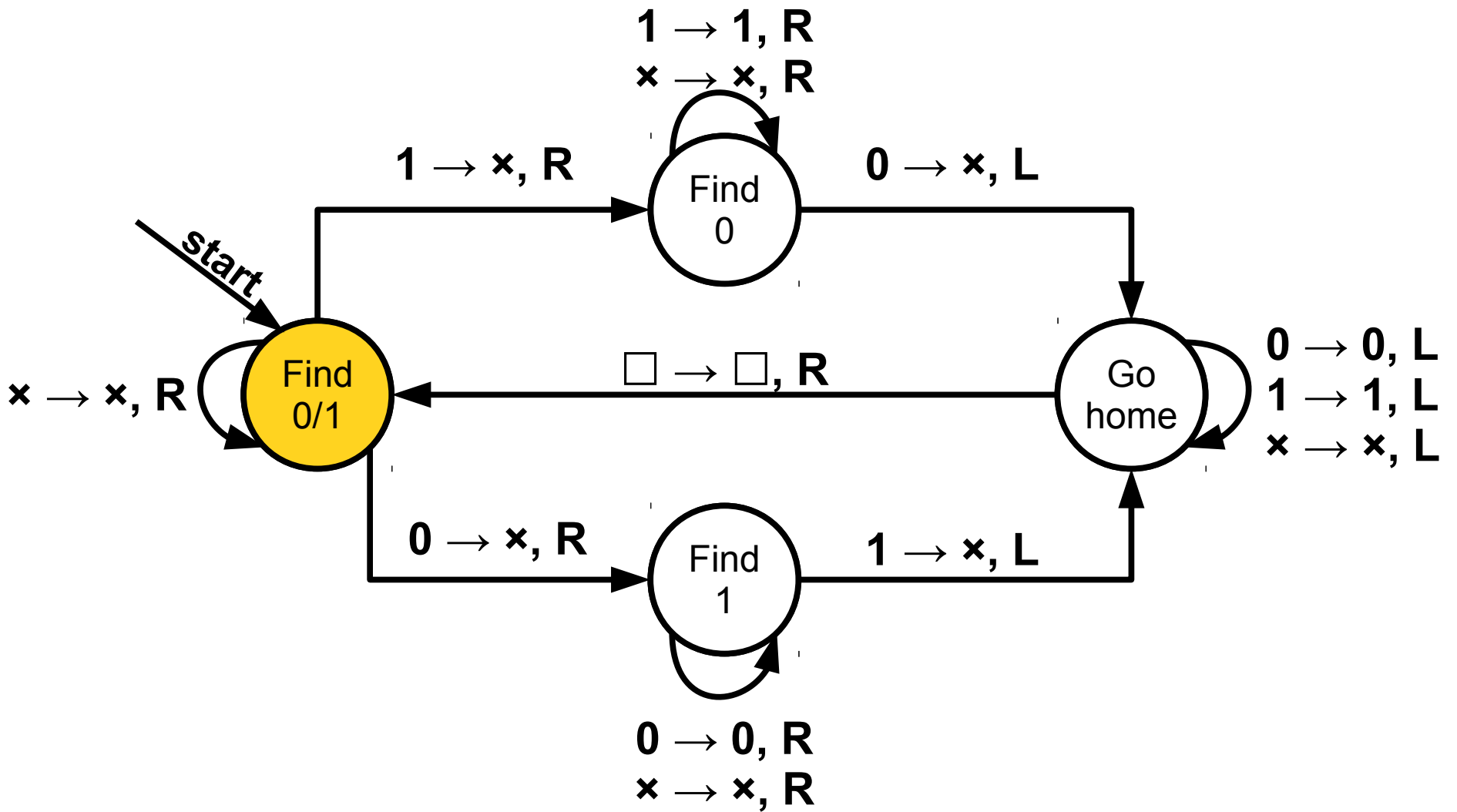


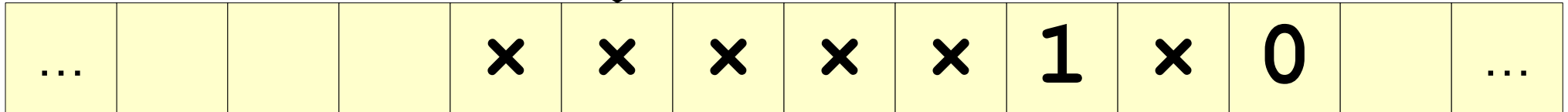
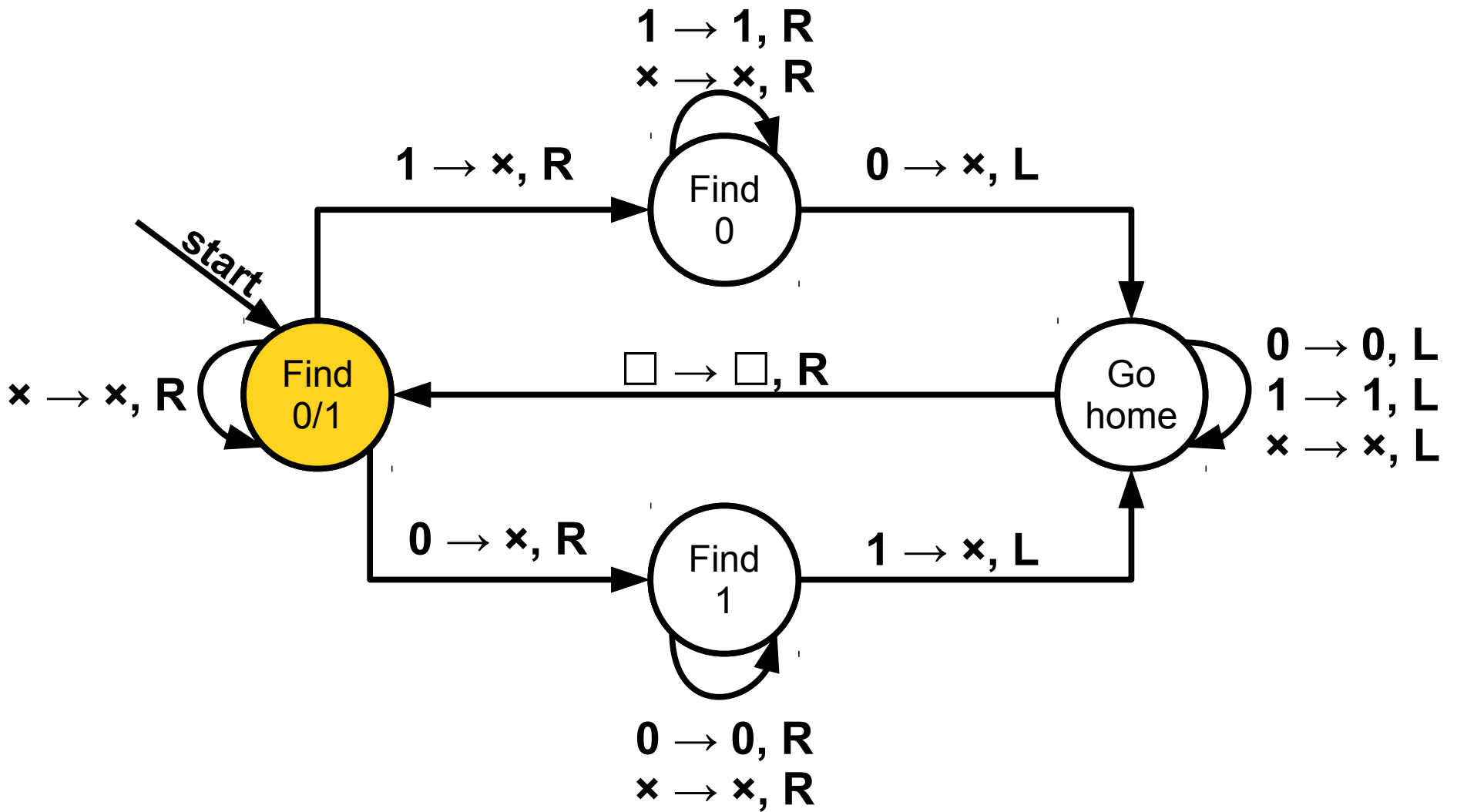


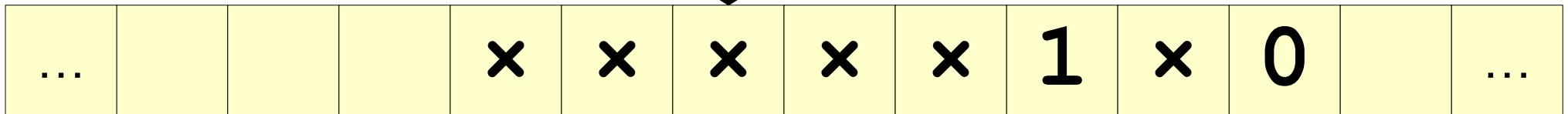
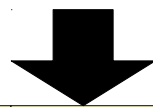
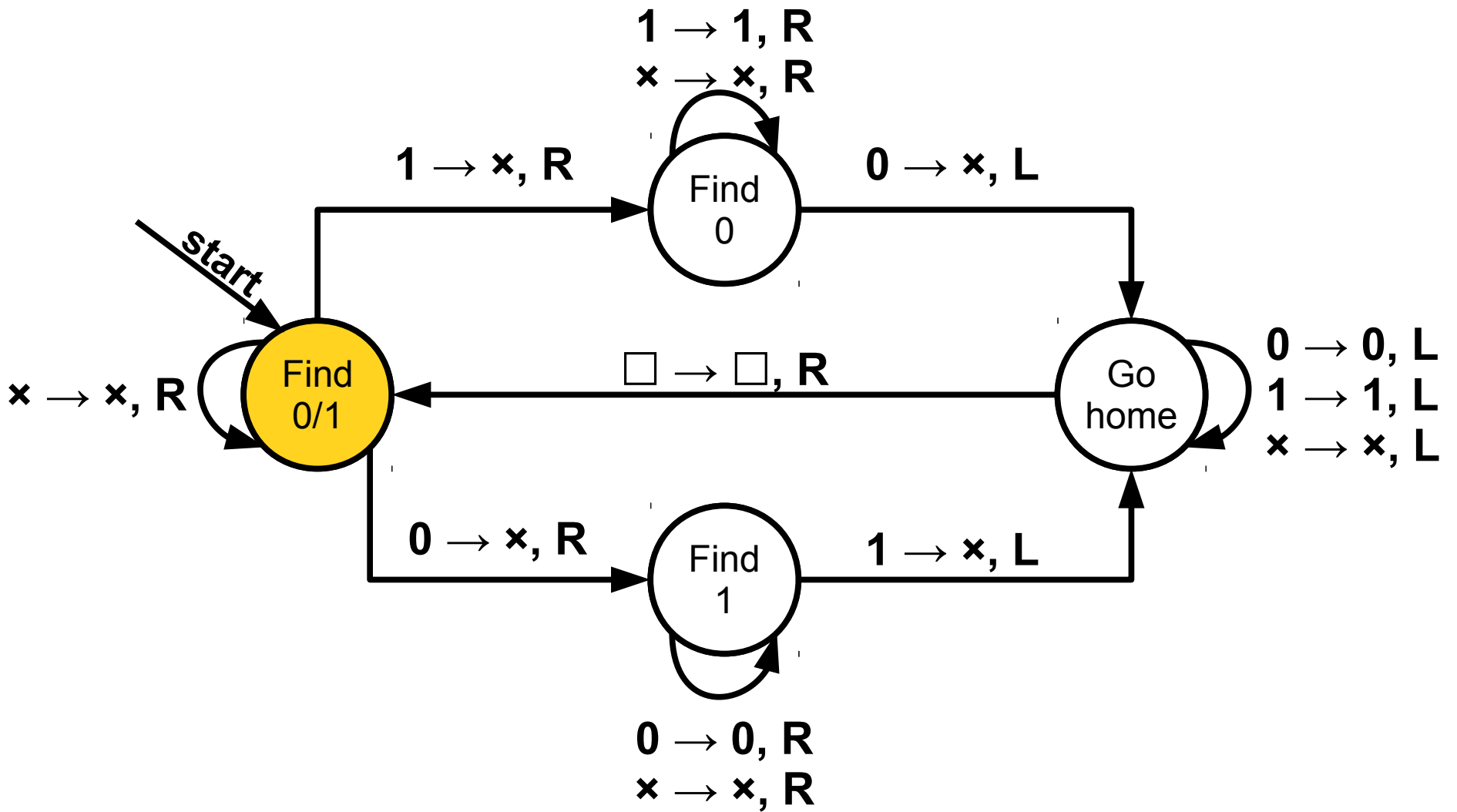


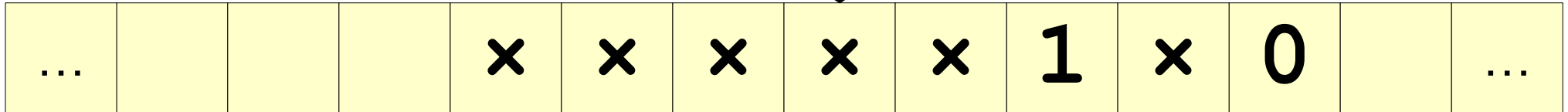
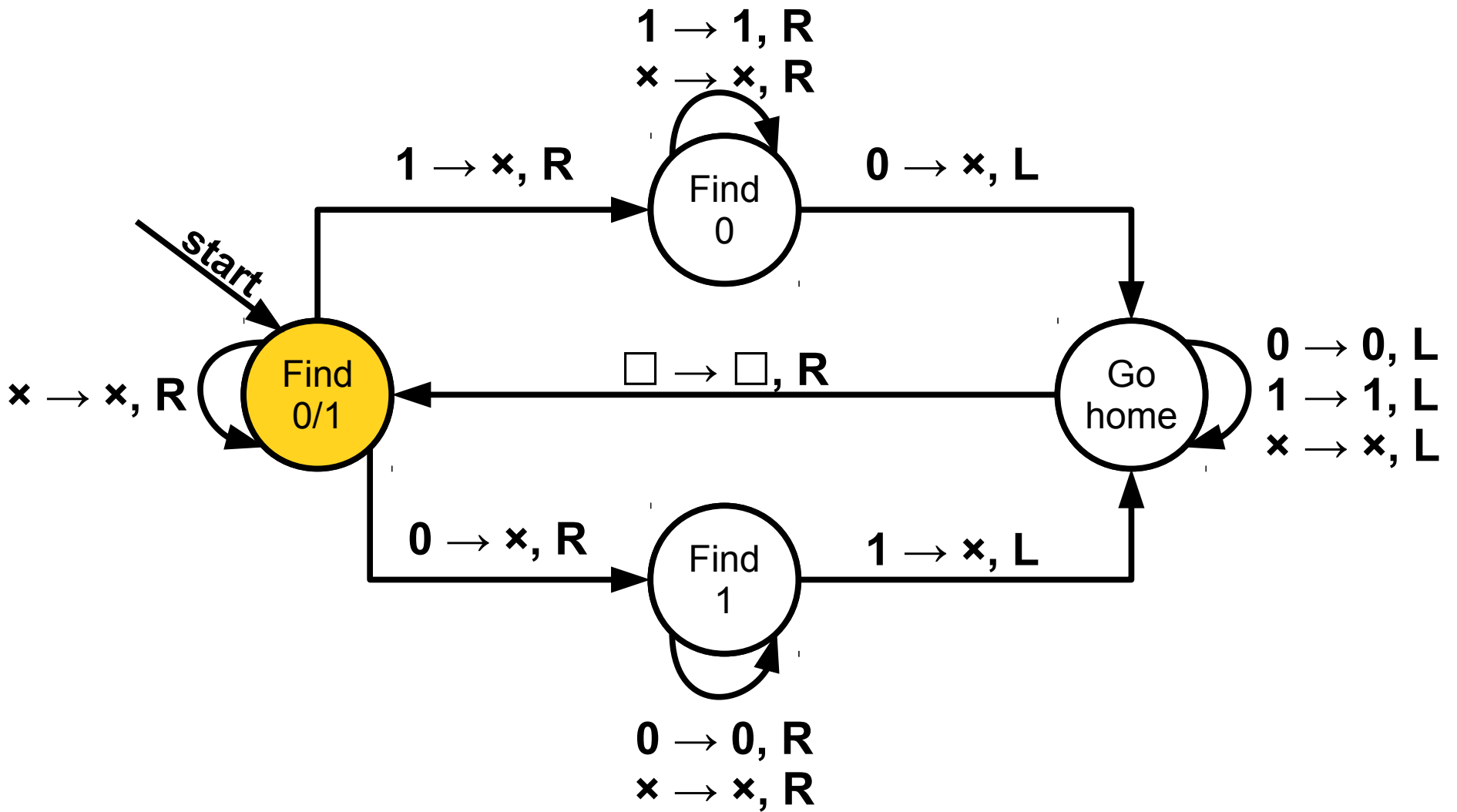


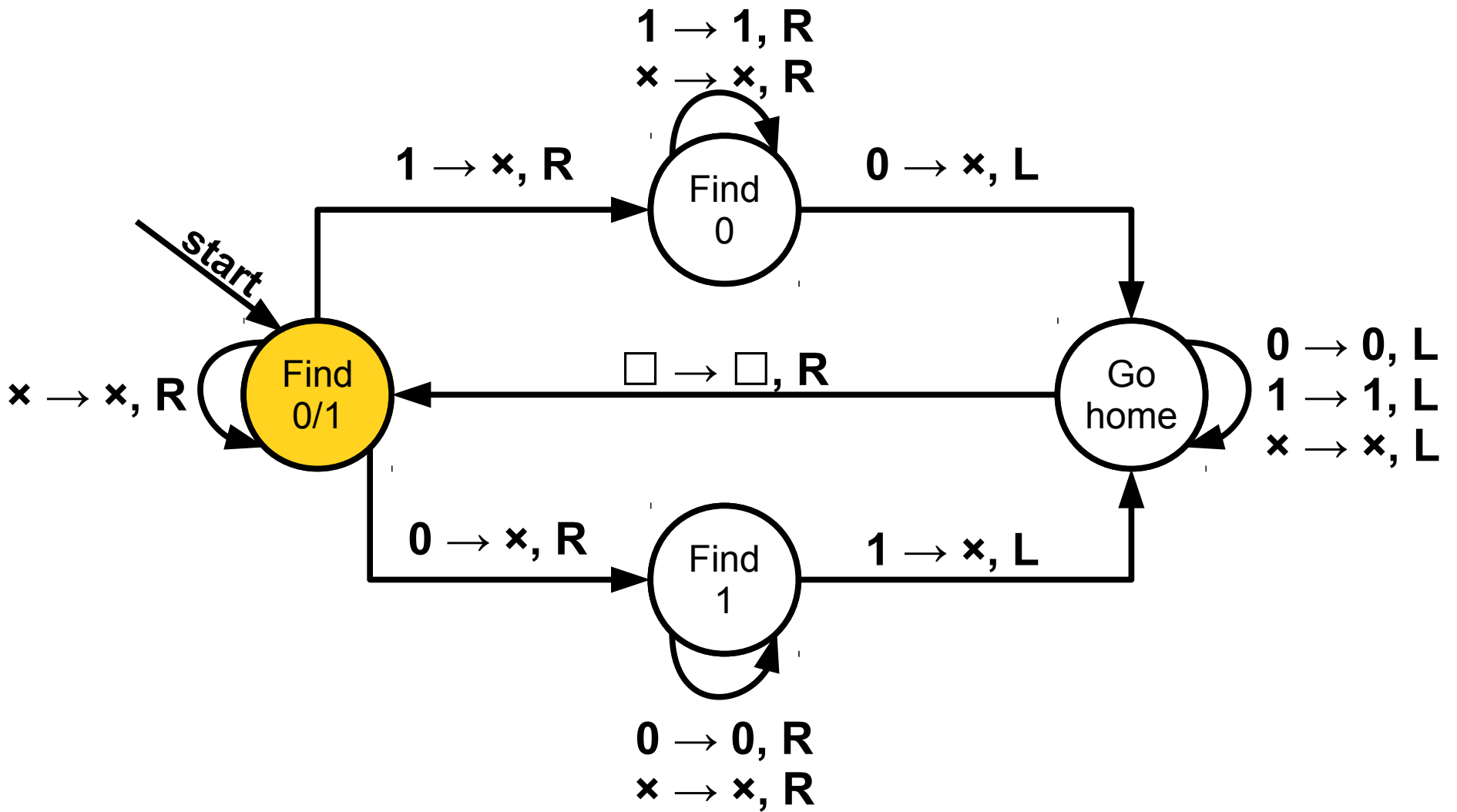


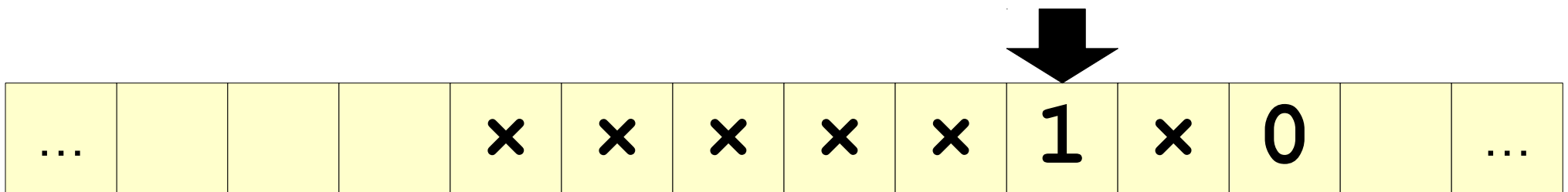
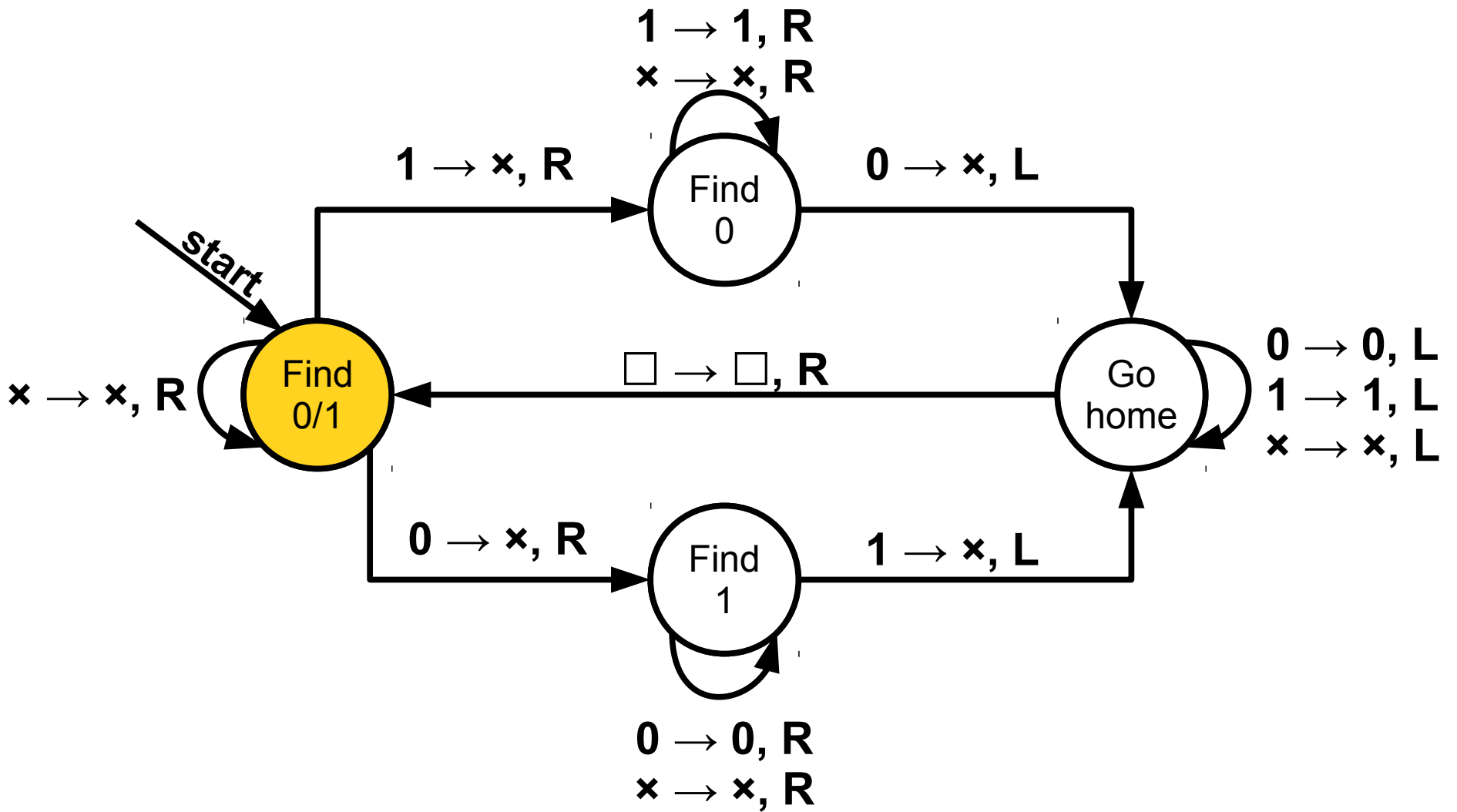


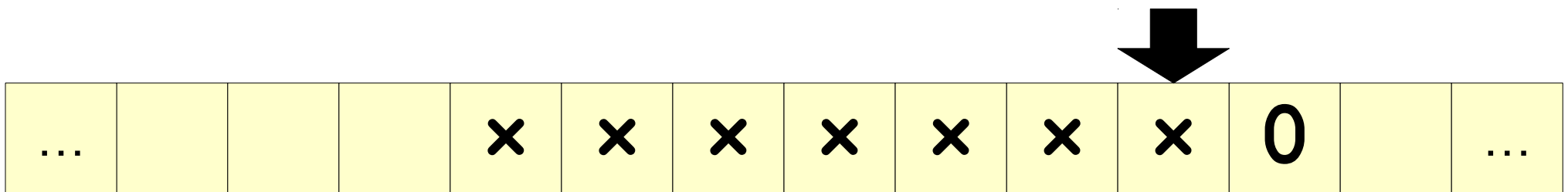
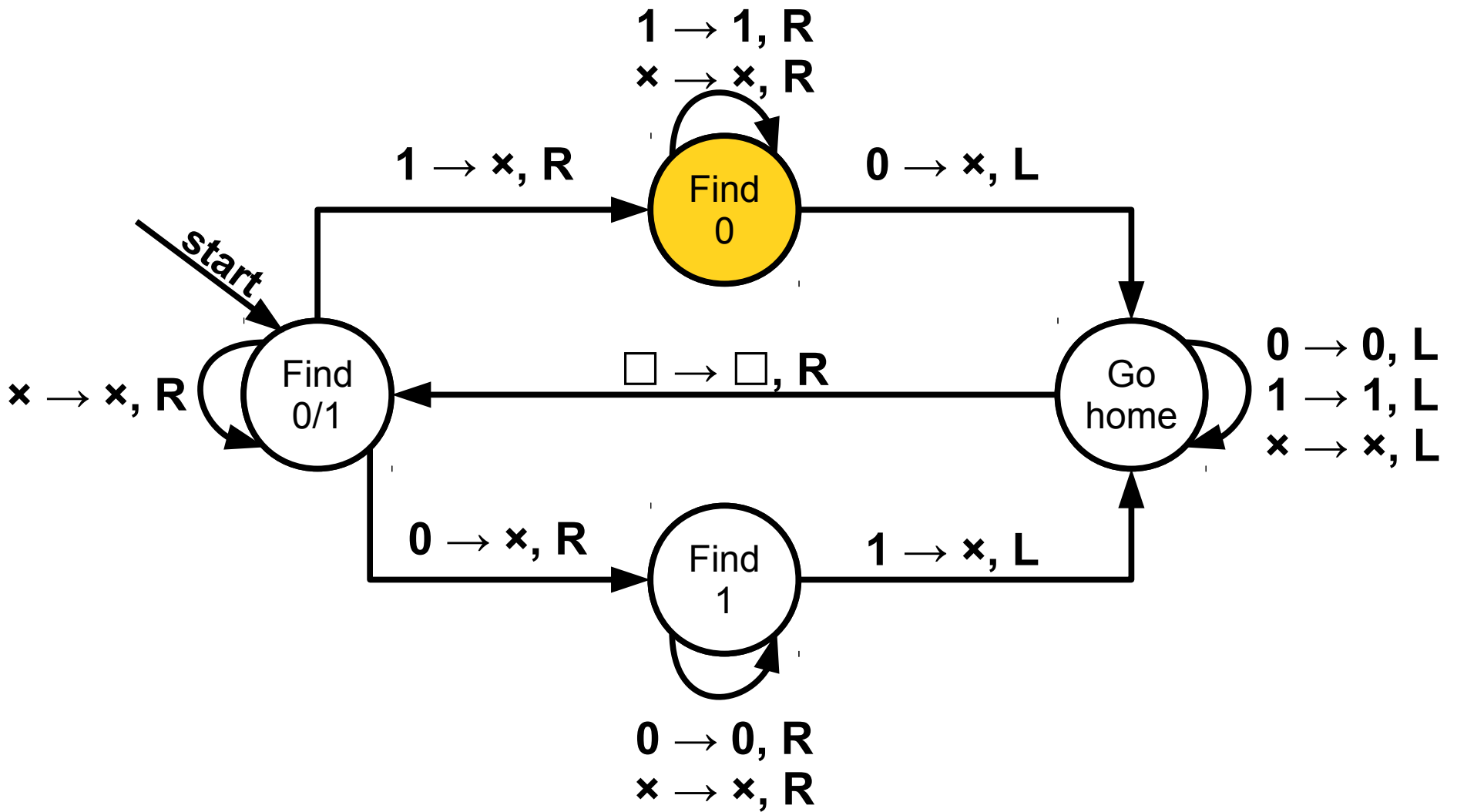




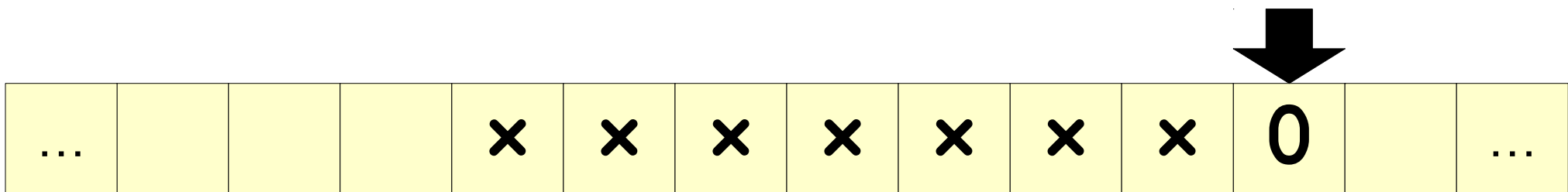
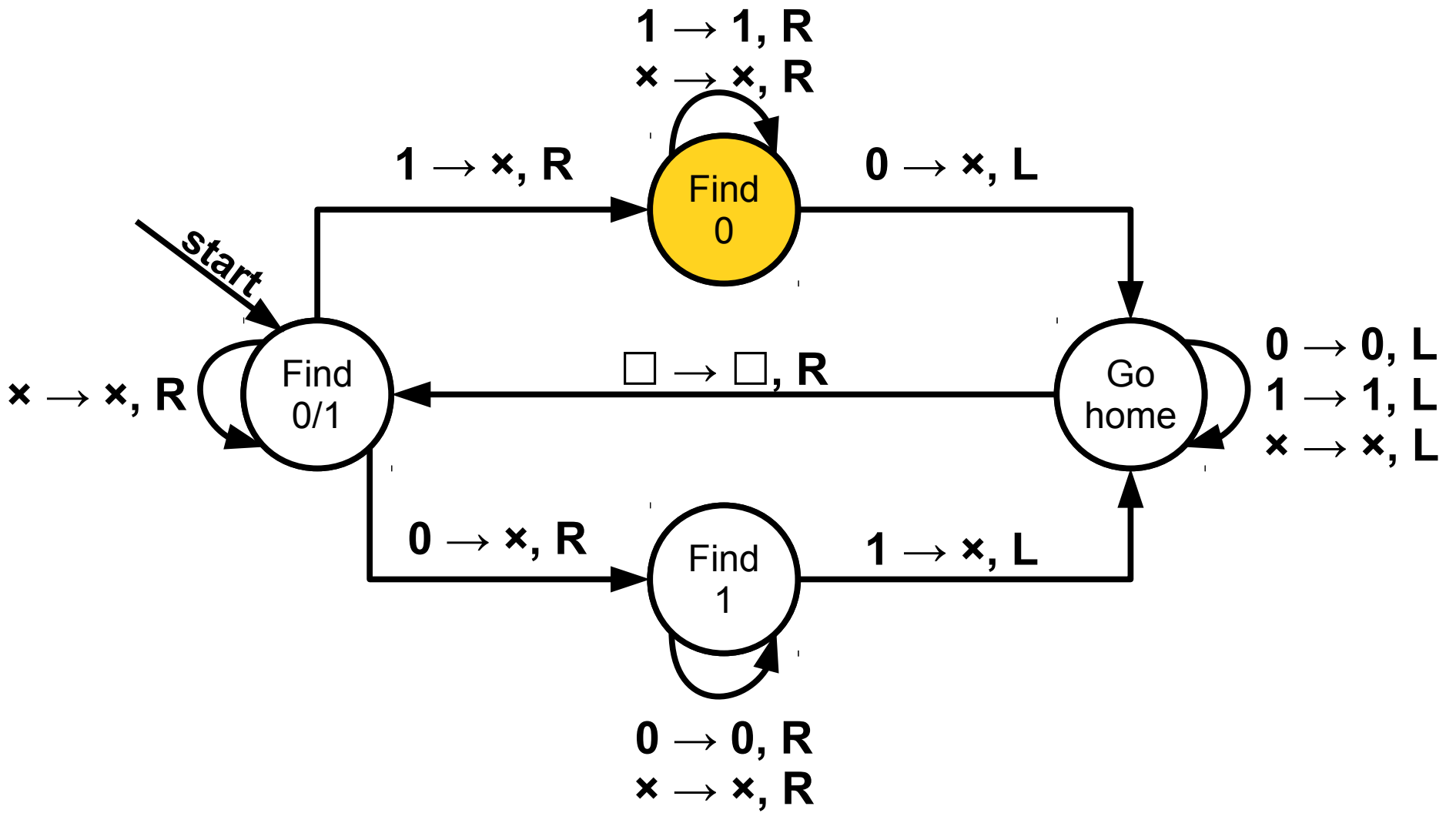


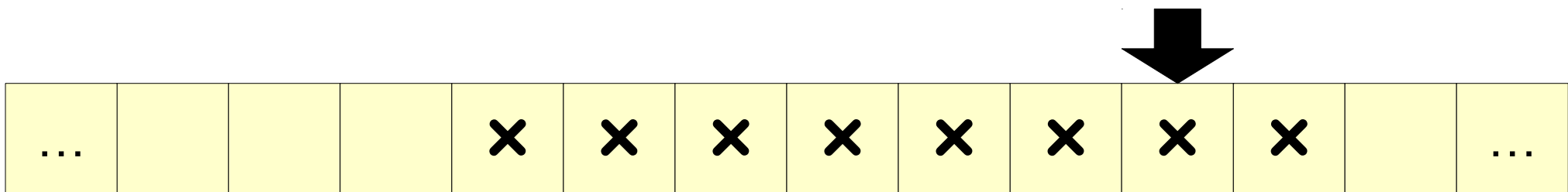
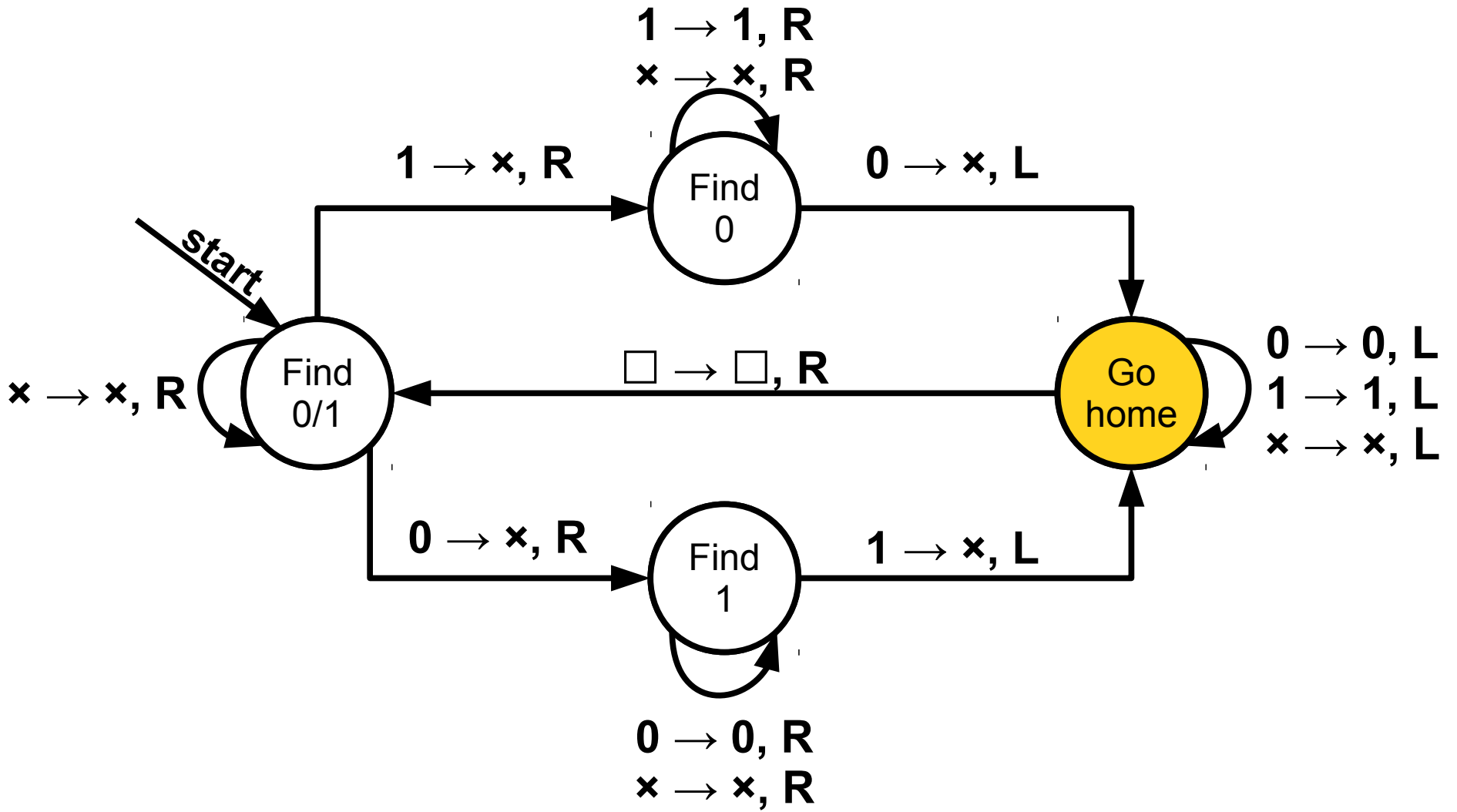




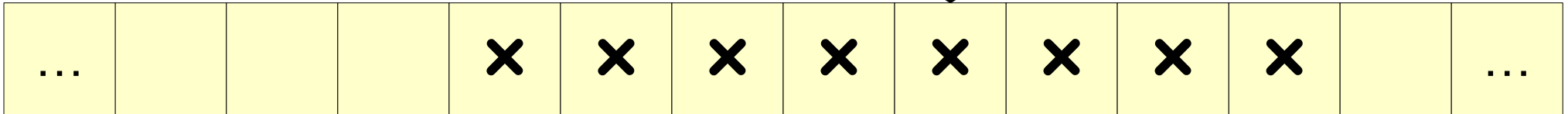
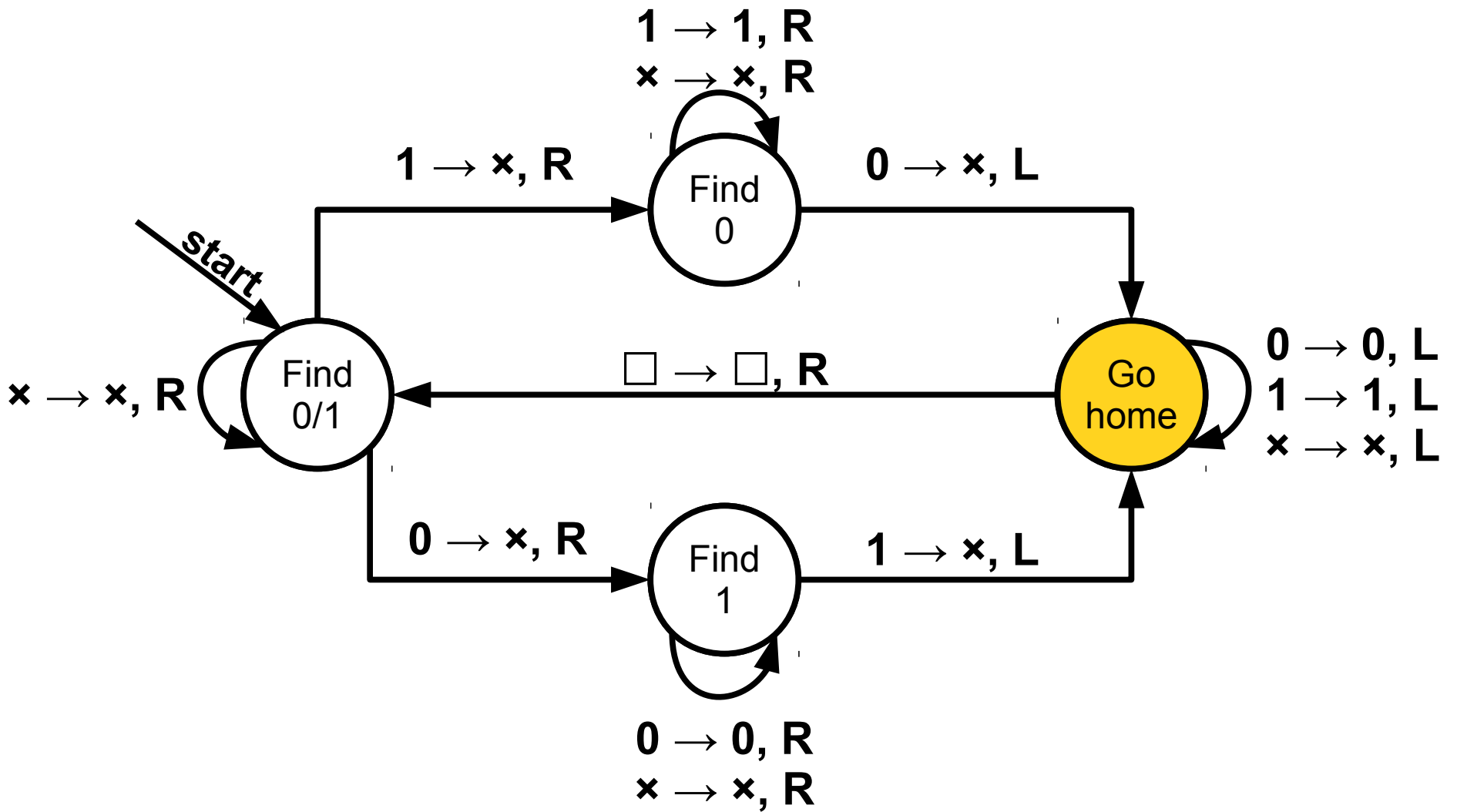














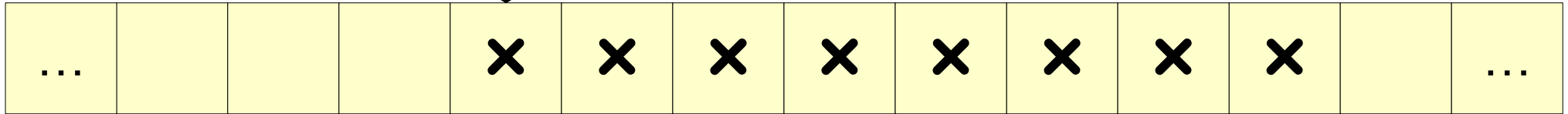
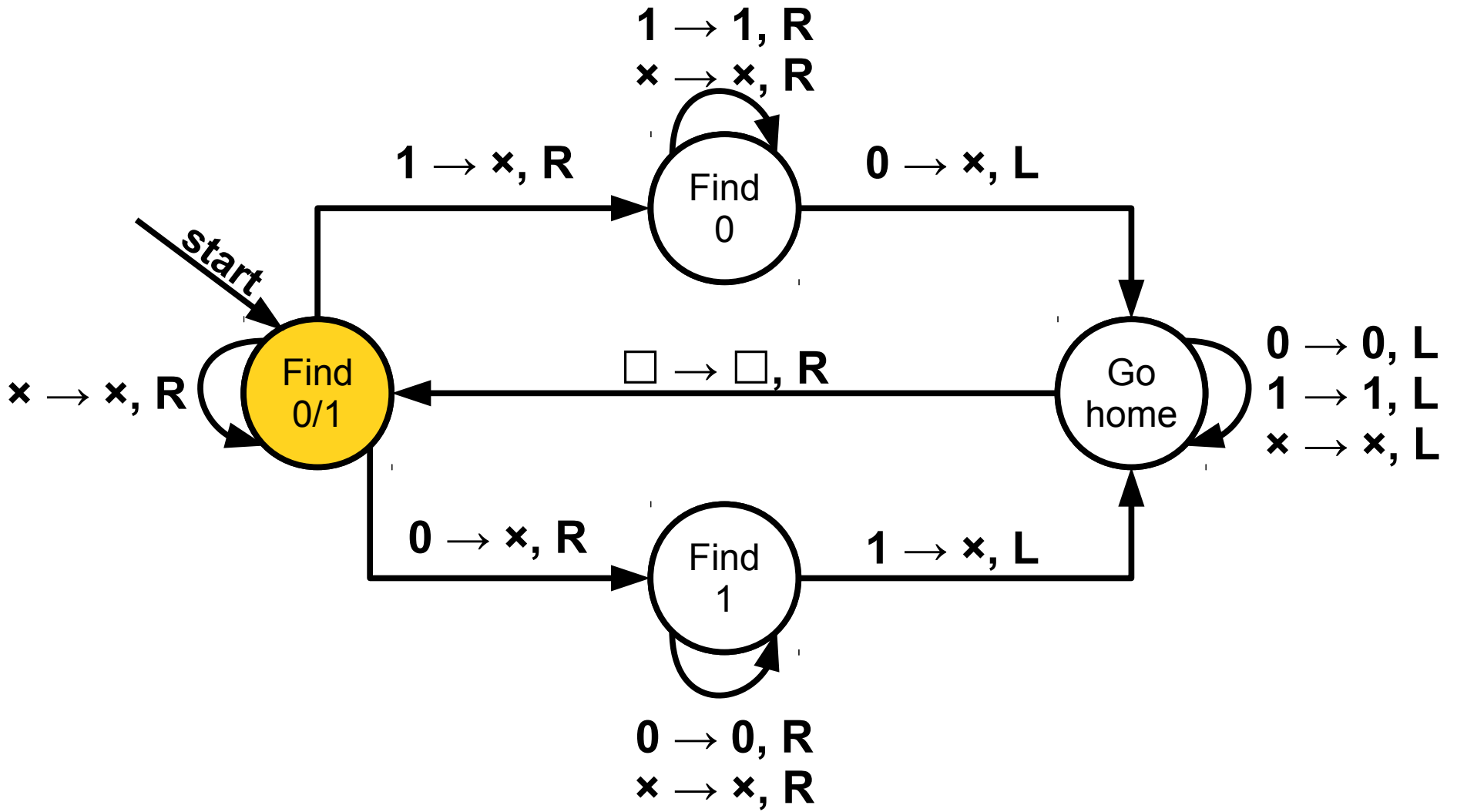


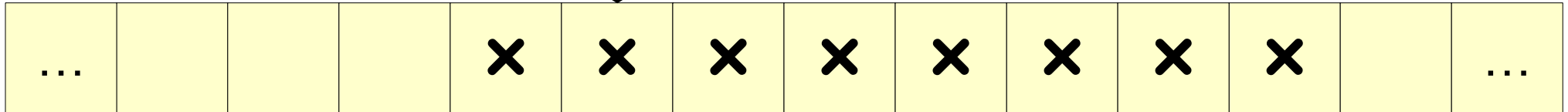
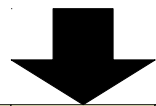
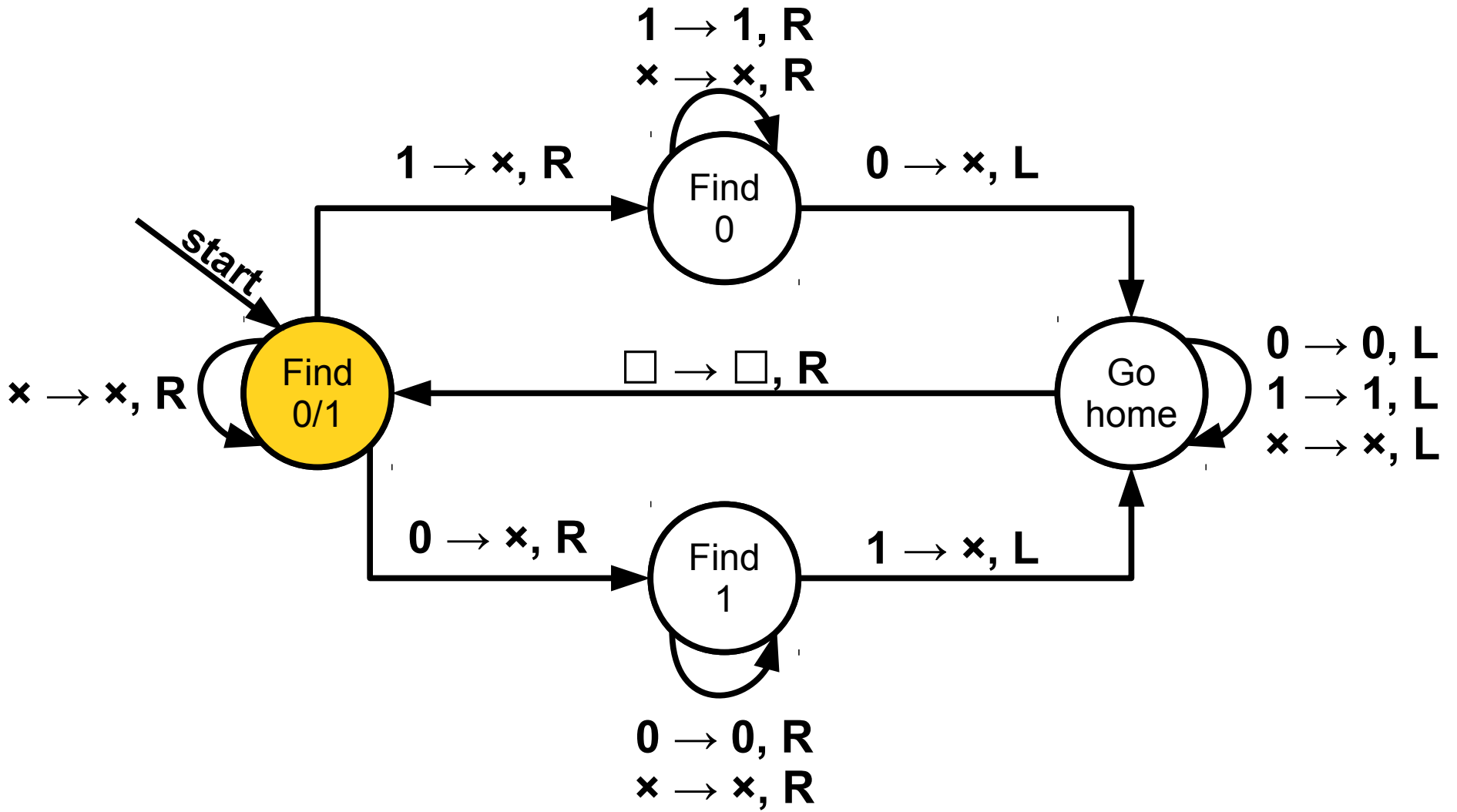


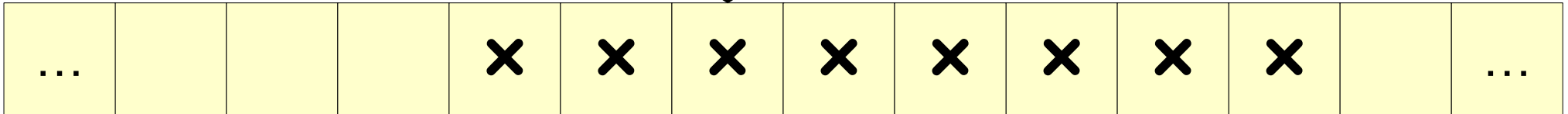
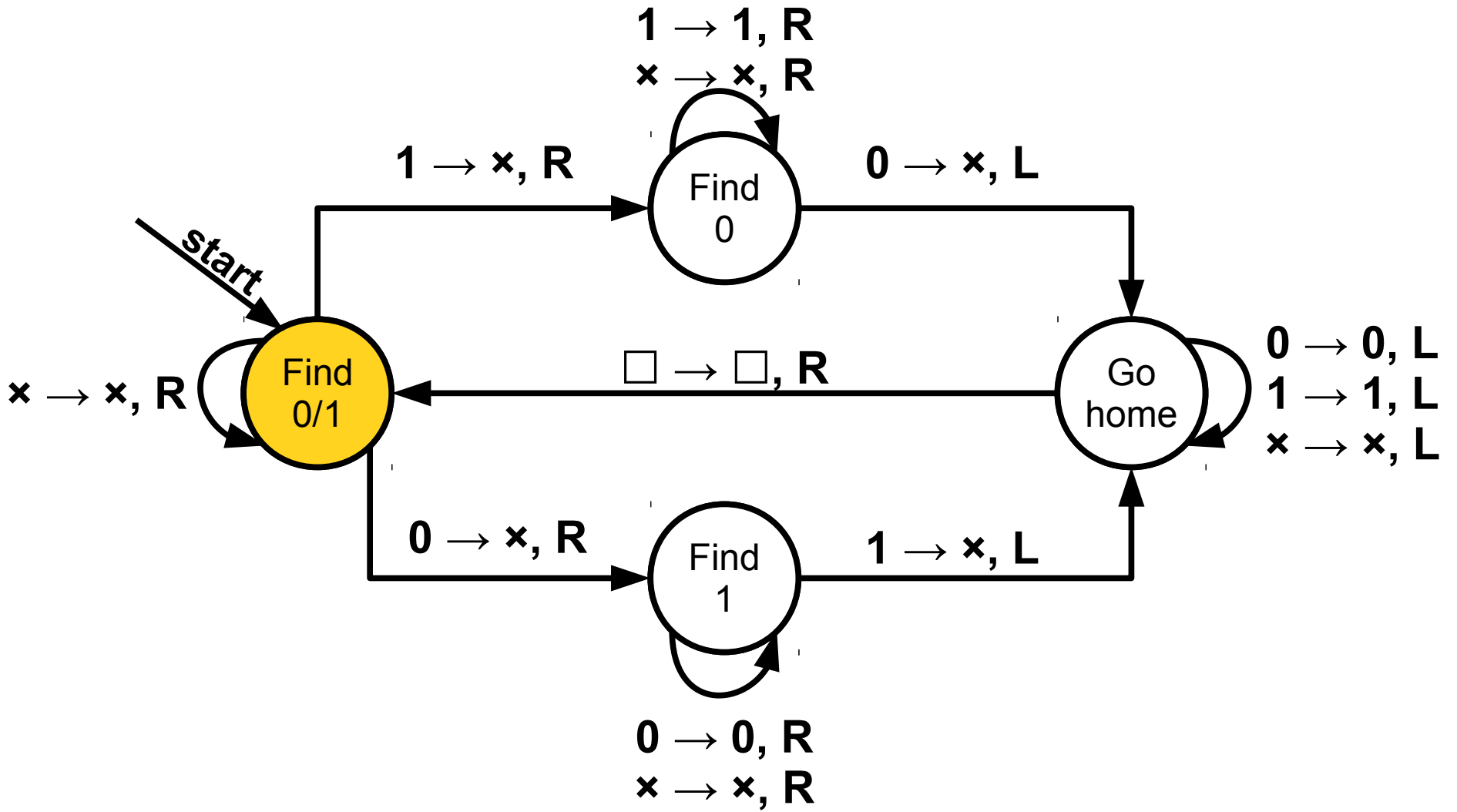




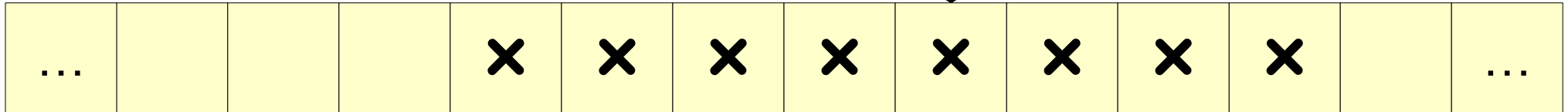
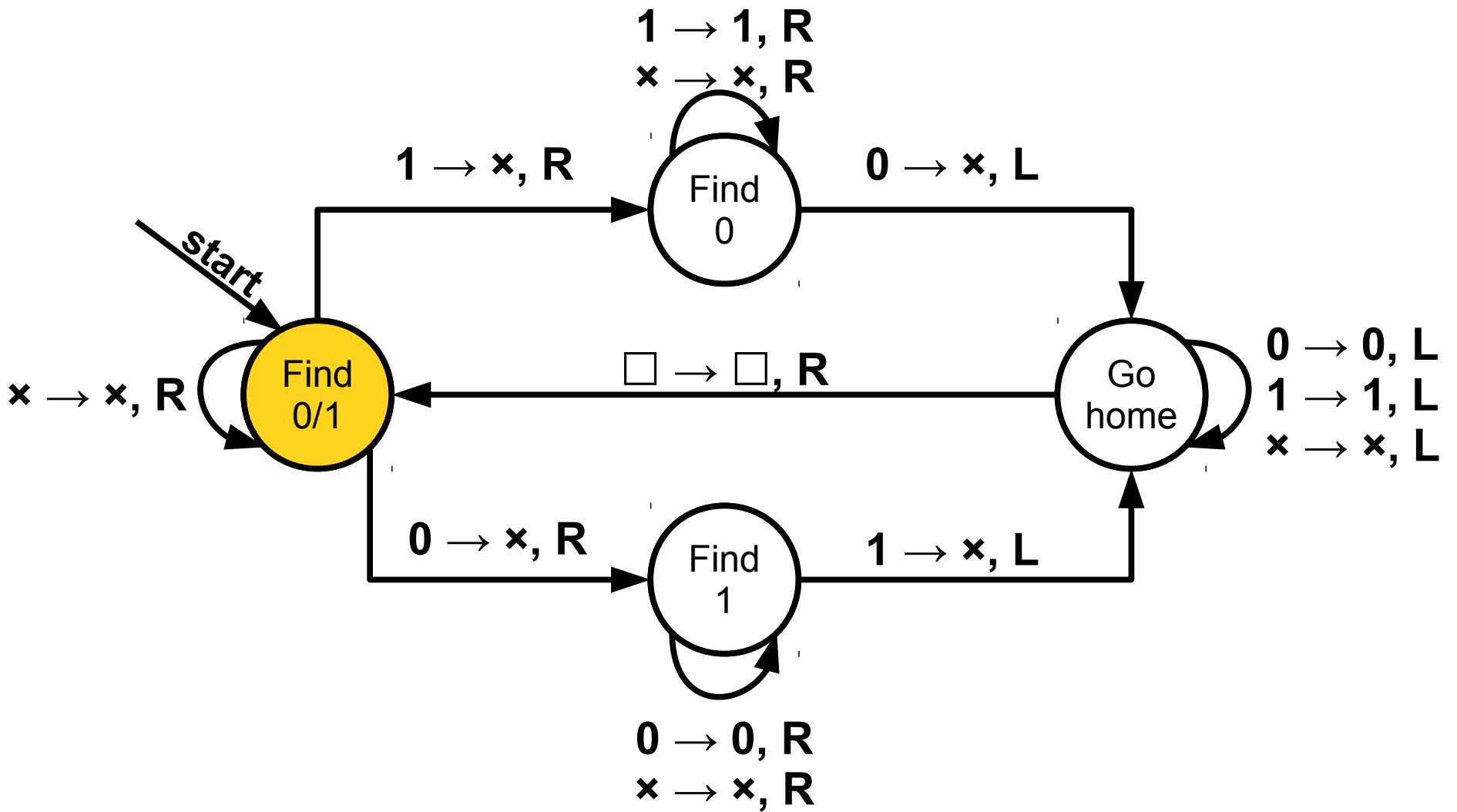








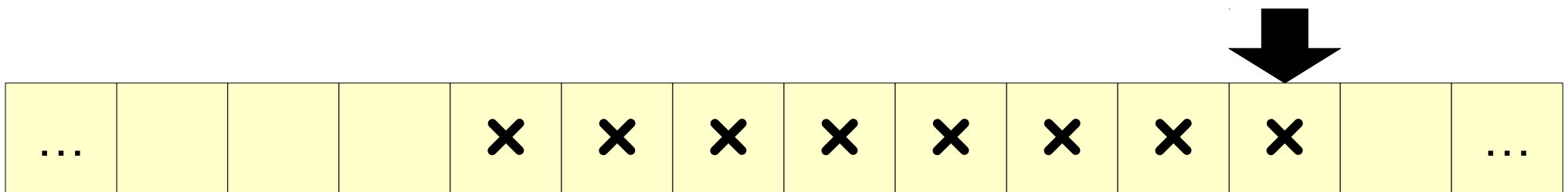
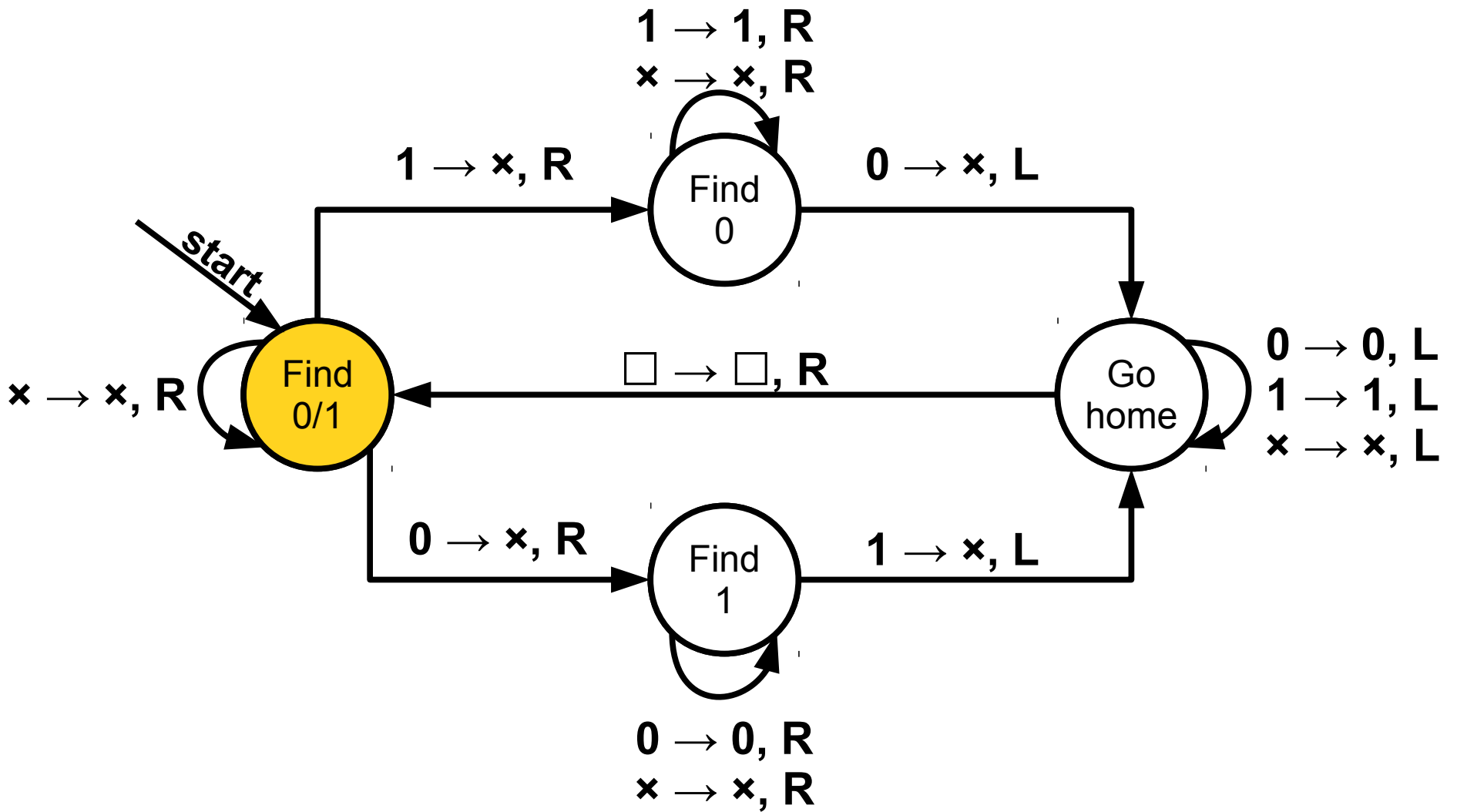


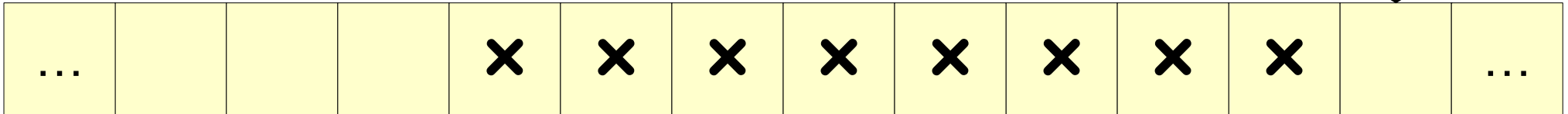
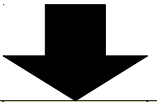
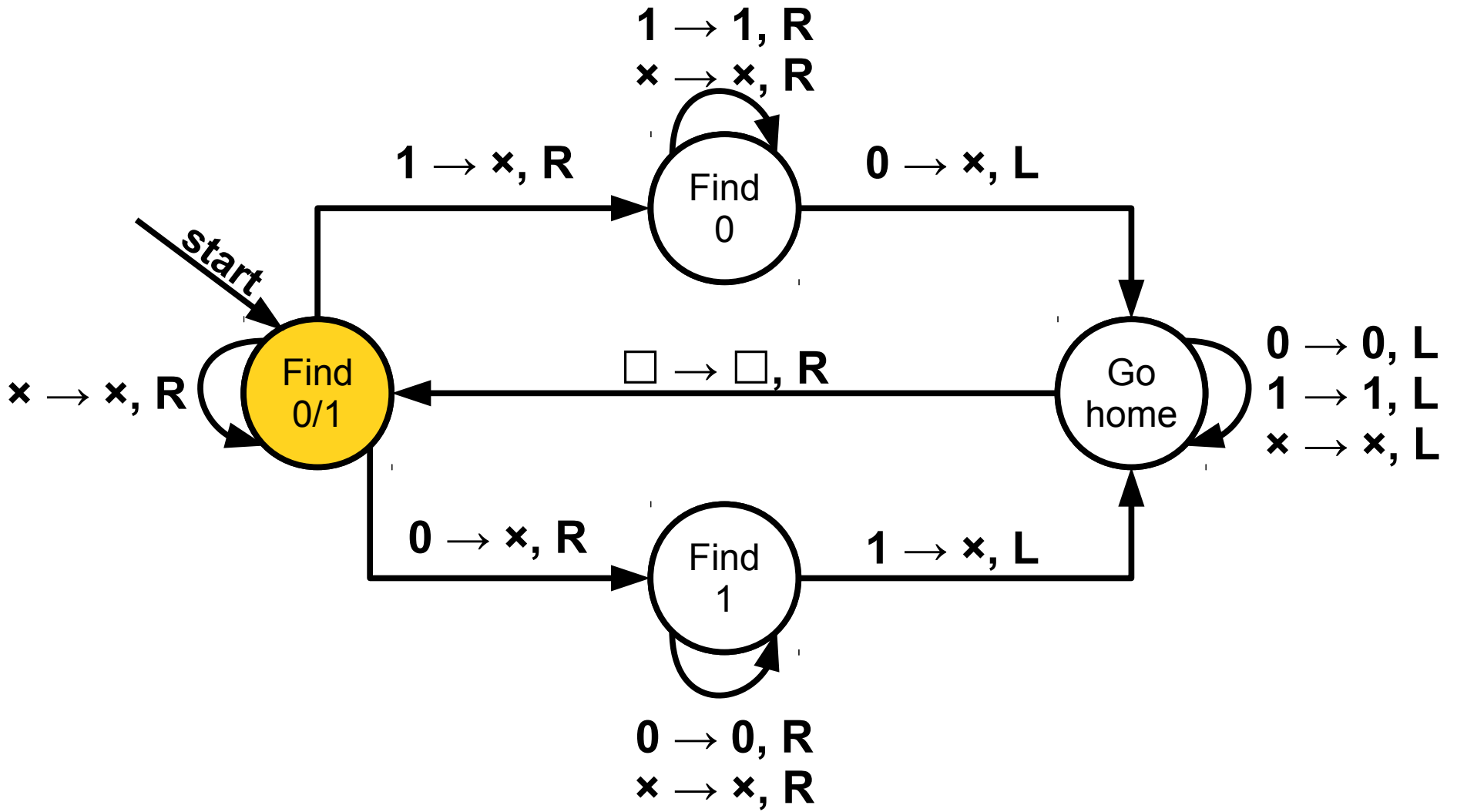


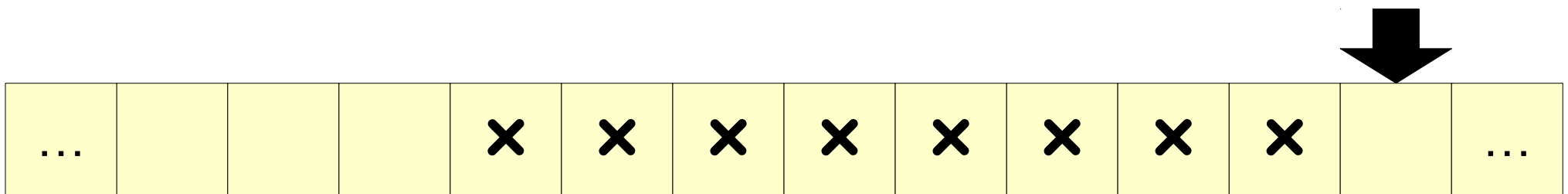
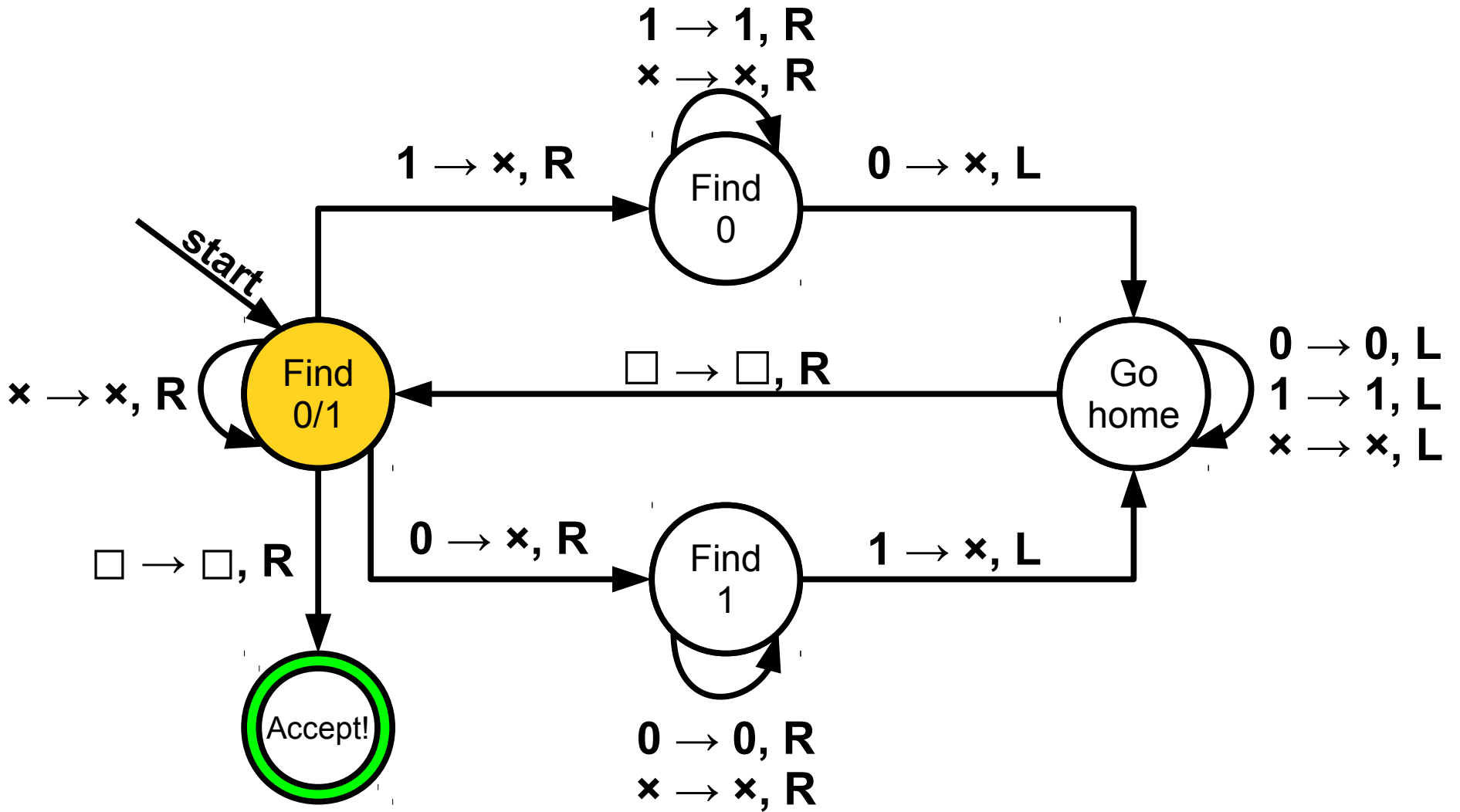


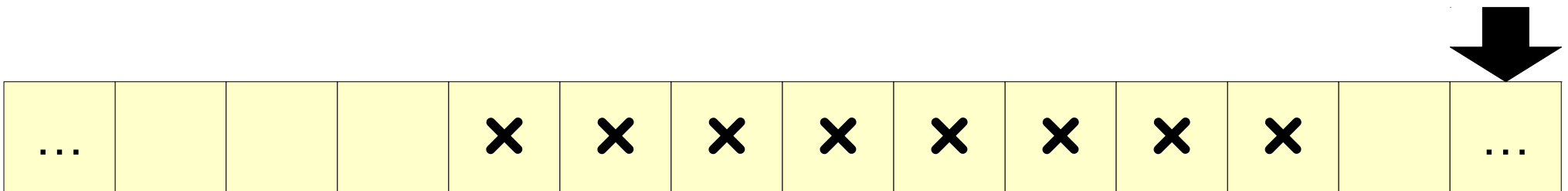
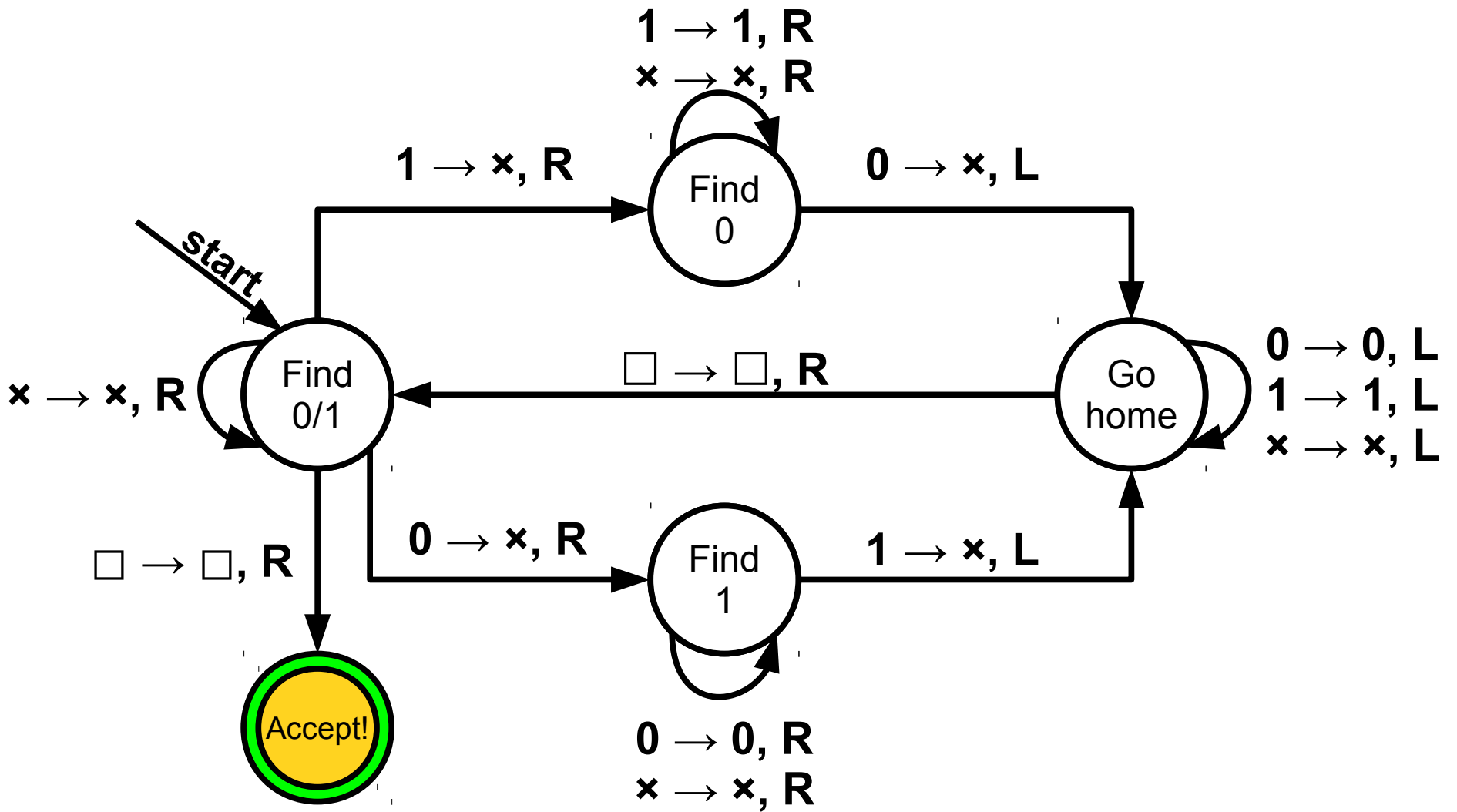


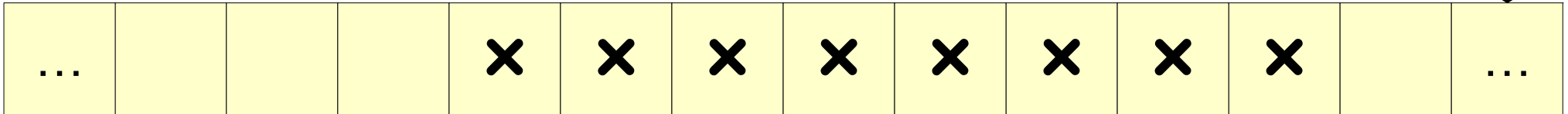
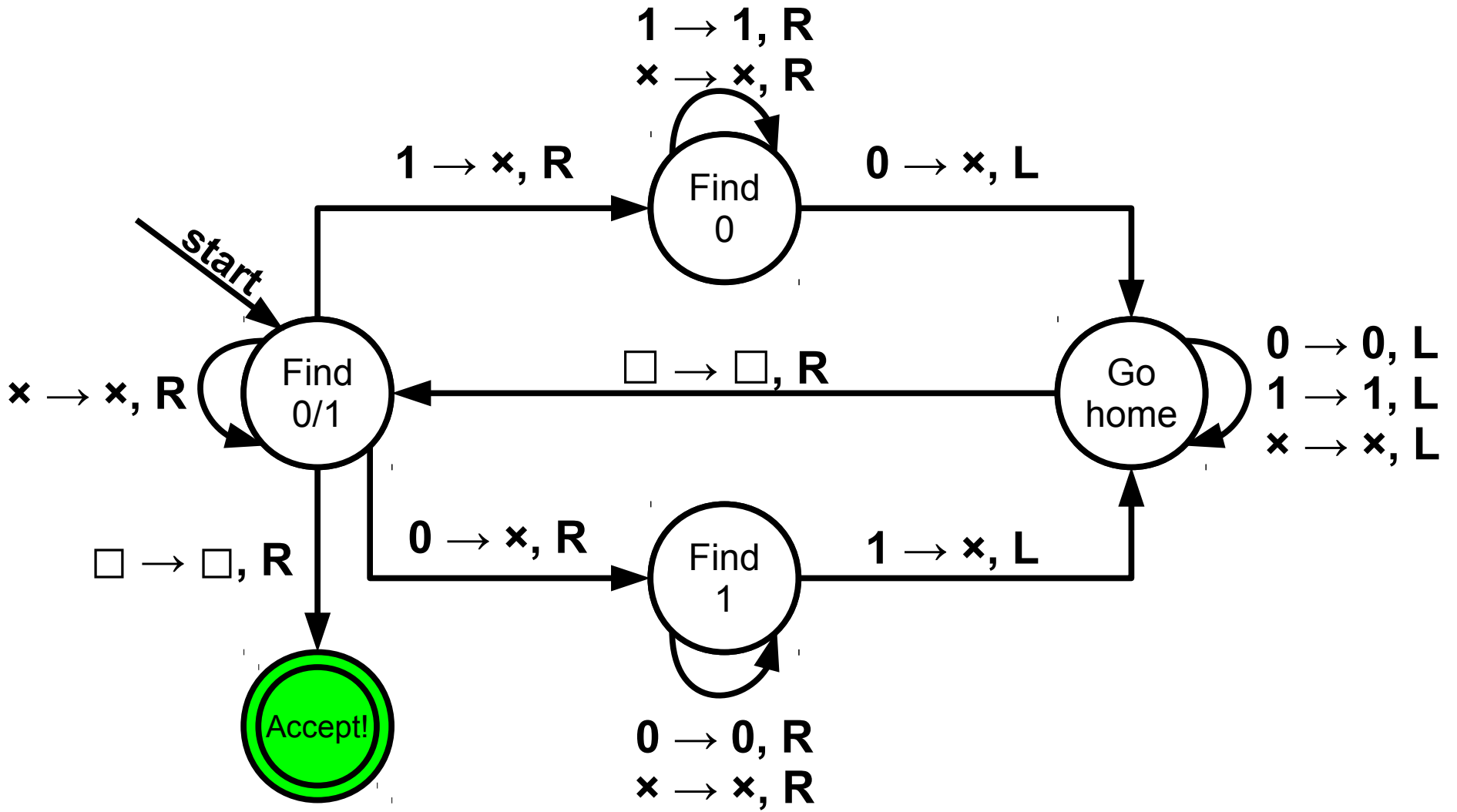


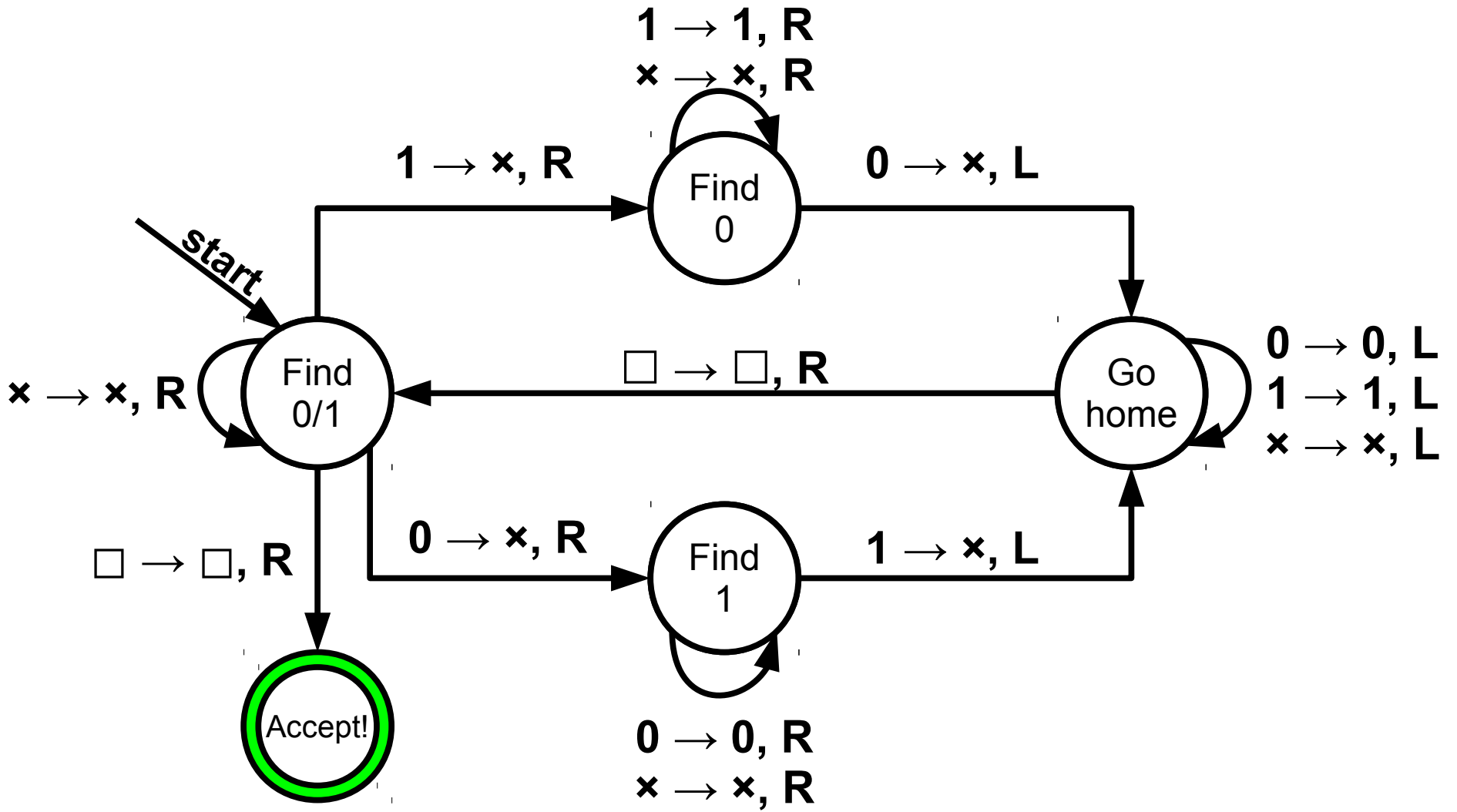


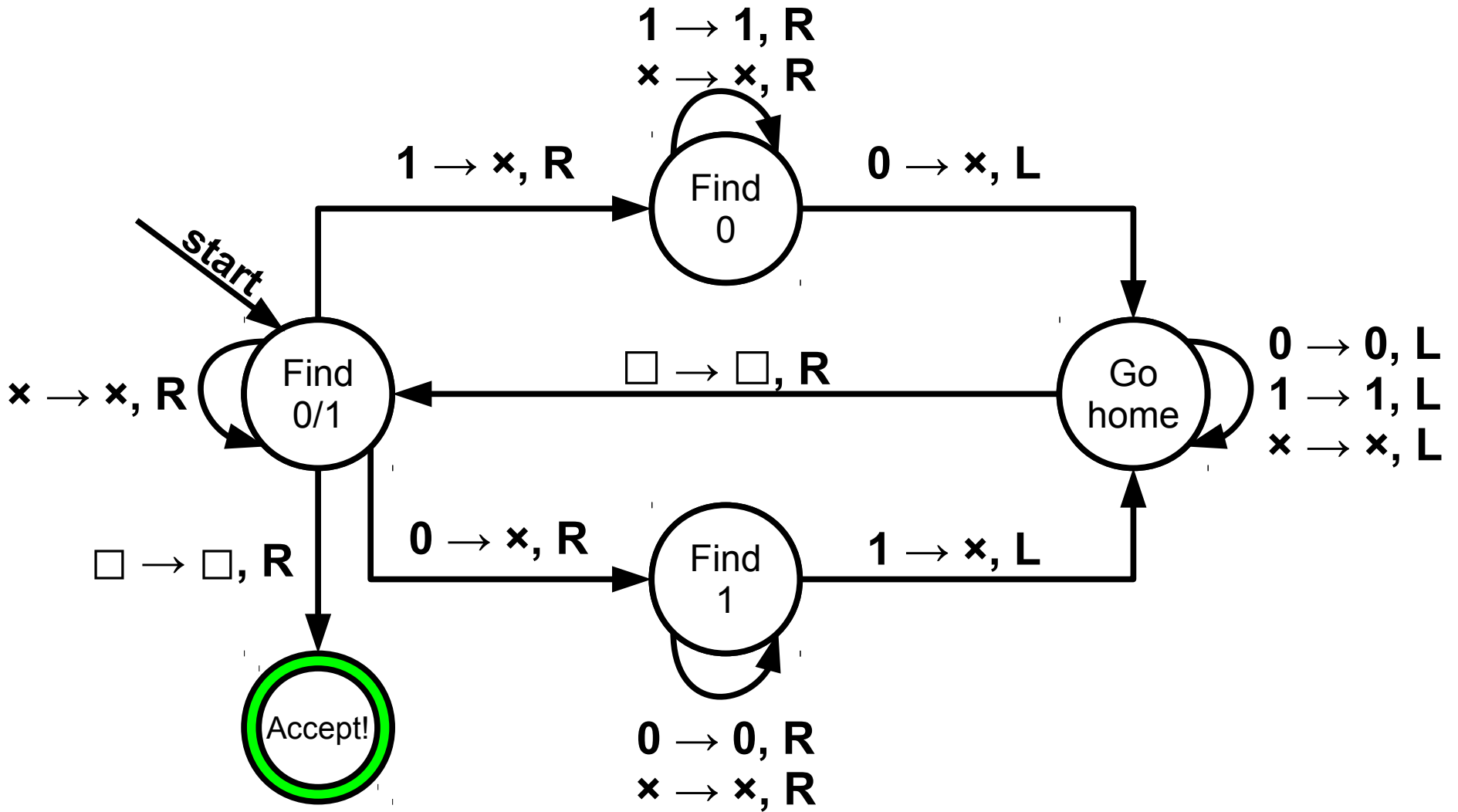




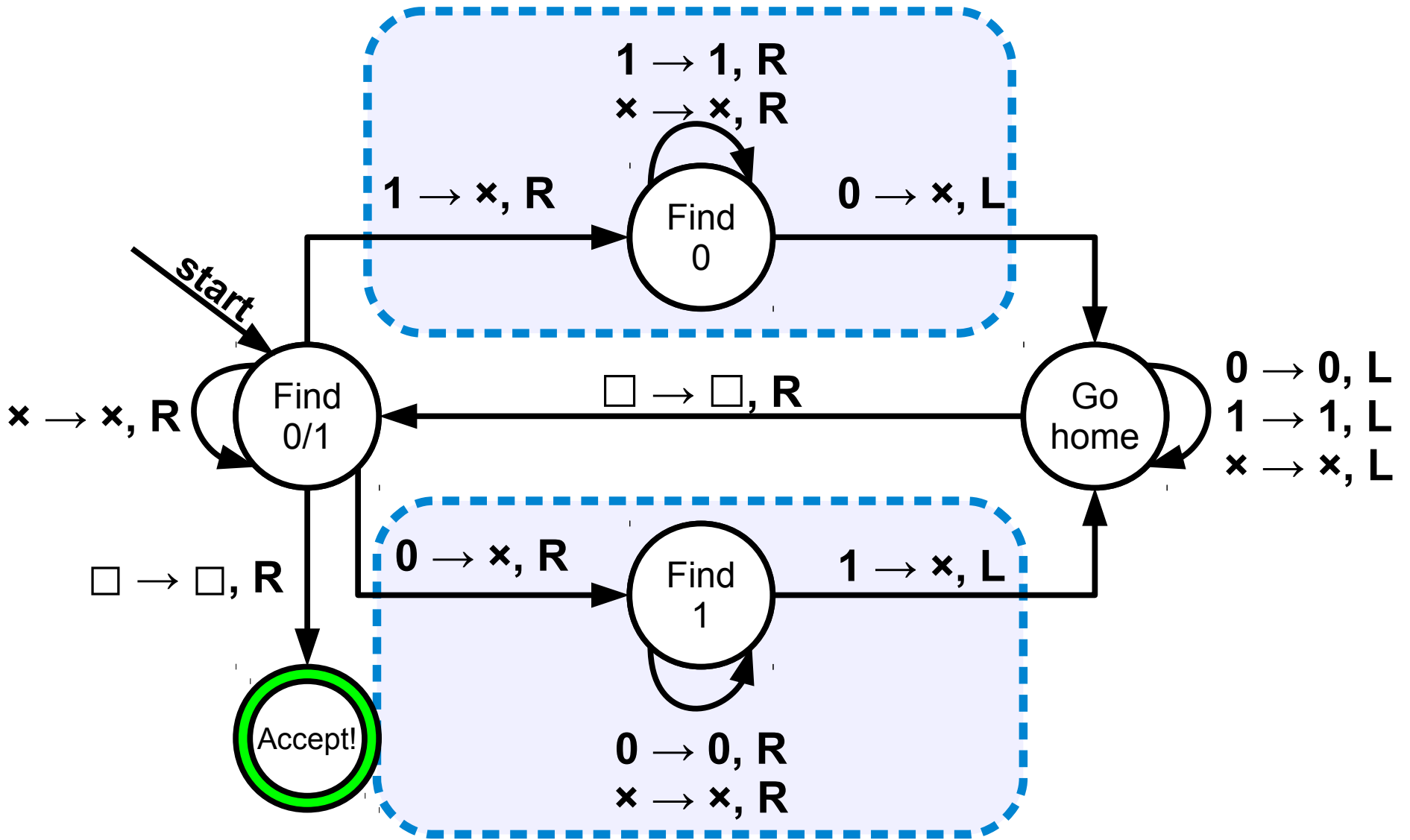








Remember that all missing transitions implicitly reject.





# Constant Storage

- Sometimes, a TM needs to remember some additional information that can't be put on the tape.
- In this case, you can use similar techniques from DFAs and introduce extra states into the TM's finite-state control.
- The finite-state control can only remember one of finitely many things, but that might be all that you need!

**Time-Out for Announcements!**

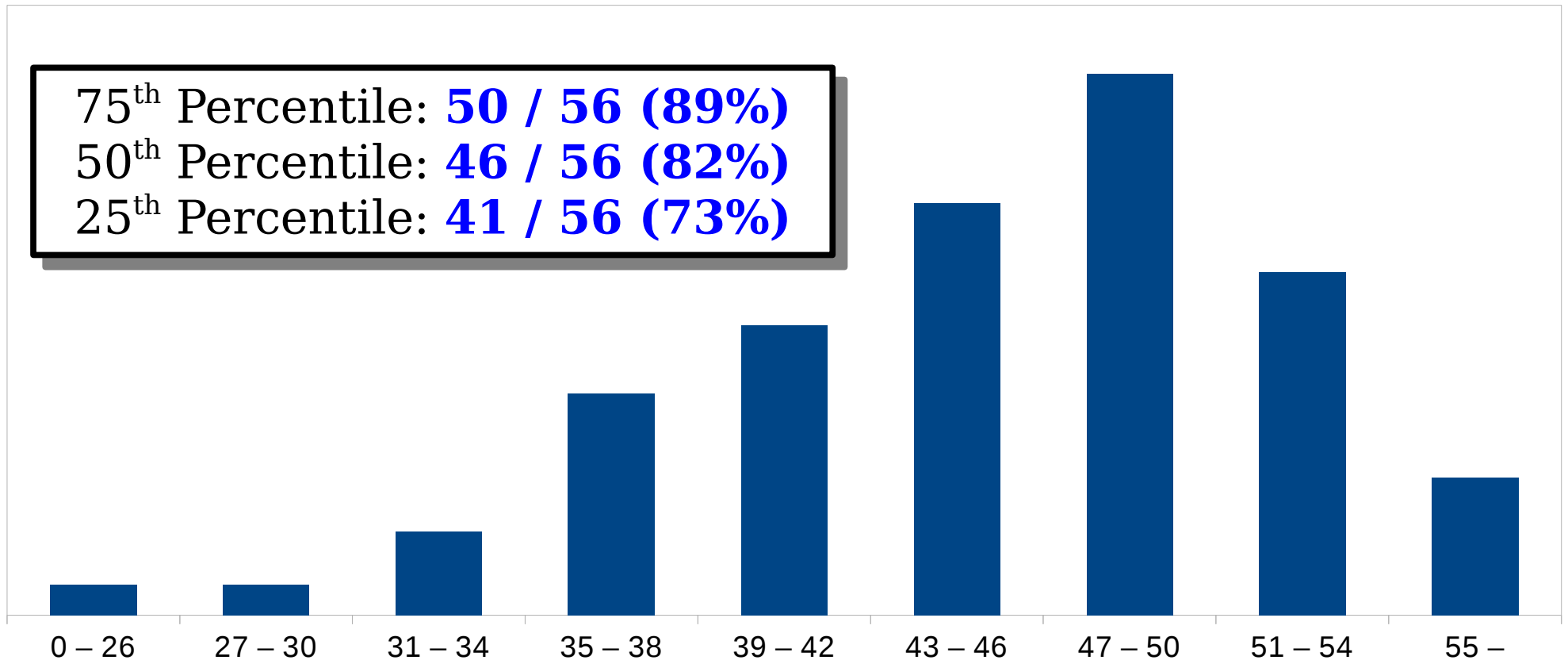
# Second Midterm Exam

- You're done with the second midterm exam! Wooahoo!
- We'll be grading the exam this weekend. Unfortunately, we will not be able to get grades back before Friday.
- Have questions? Feel free to ask in office hours or on Piazza!

# Problem Set Seven

- Problem Set Seven is due this Friday at 2:30PM.
- As always, if you have questions, feel free to stop by office hours or ask on Piazza!

# Problem Set Six Scores



Back to CS103!

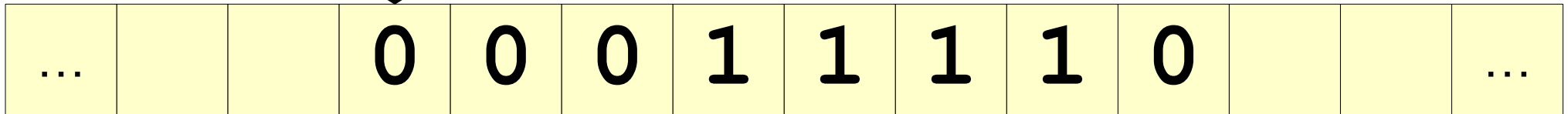
# Another TM Design

- We just designed a TM for this language over  $\Sigma = \{0, 1\}$ :

$$L = \{ w \in \Sigma^* \mid w \text{ has the same number of } 0\text{s and } 1\text{s} \}$$

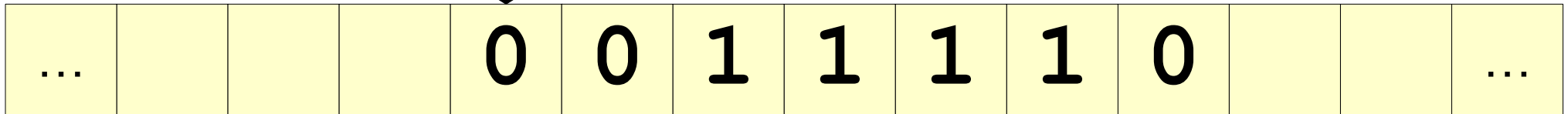
- Let's do a quick review of how it worked.

# A Leap of Faith

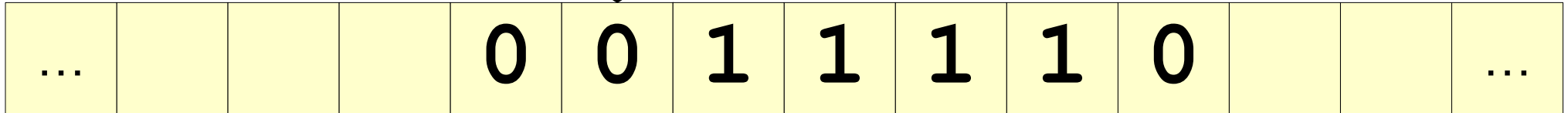




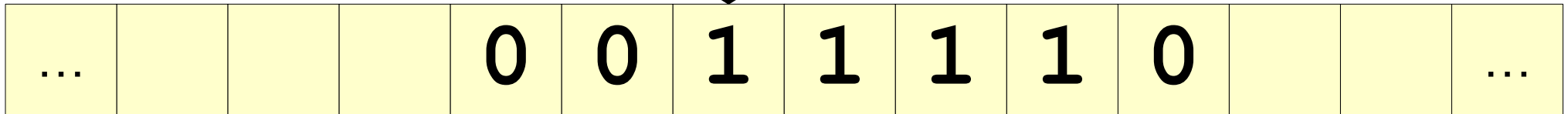
# A Leap of Faith



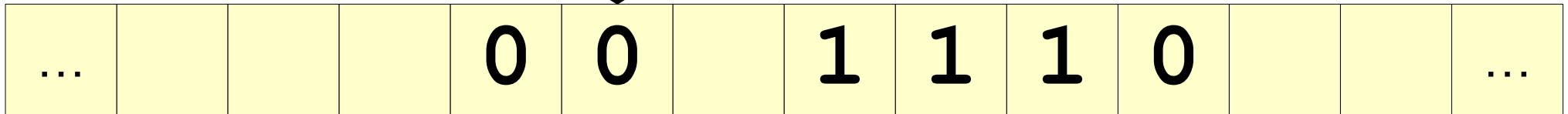
# A Leap of Faith



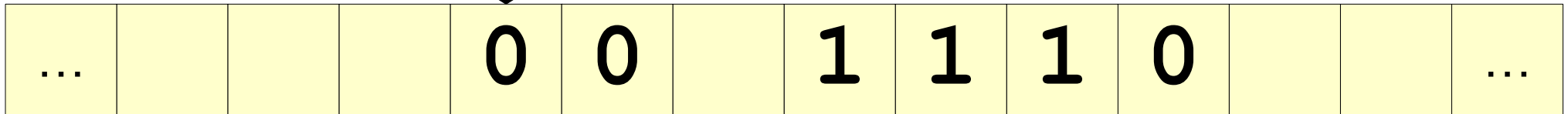
# A Leap of Faith



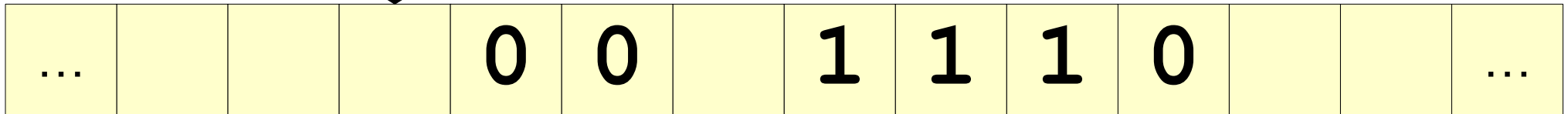
# A Leap of Faith



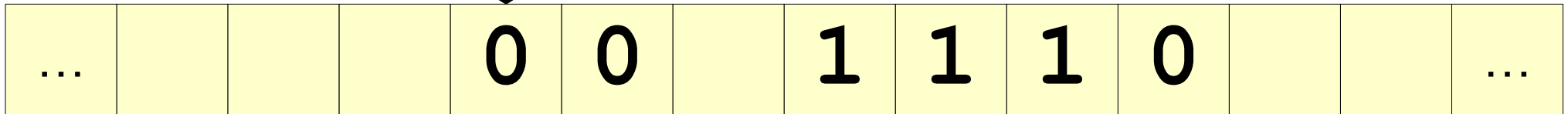
# A Leap of Faith



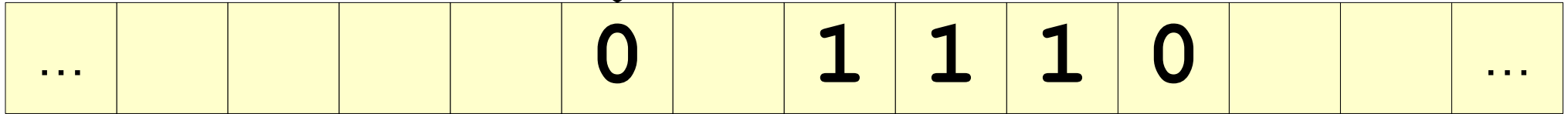
# A Leap of Faith



# A Leap of Faith

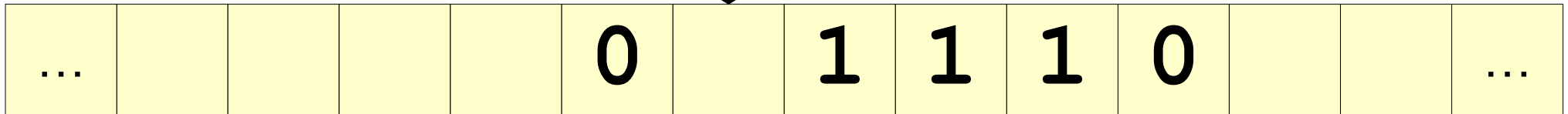


# A Leap of Faith

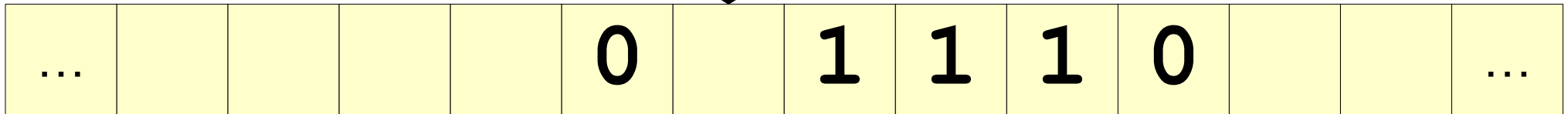




# A Leap of Faith

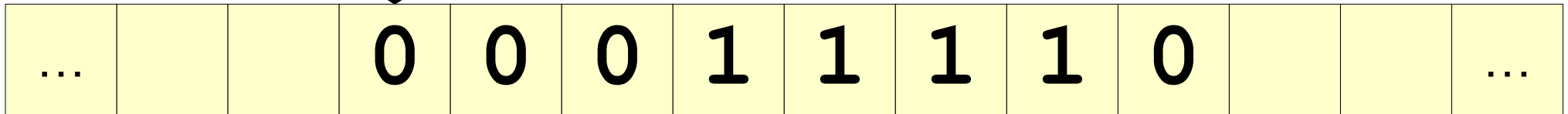


# A Leap of Faith

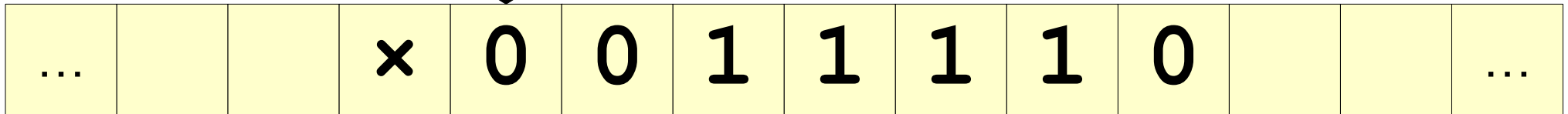


How do we know that  
this blank isn't one of  
the infinitely many  
blanks after our input  
string?

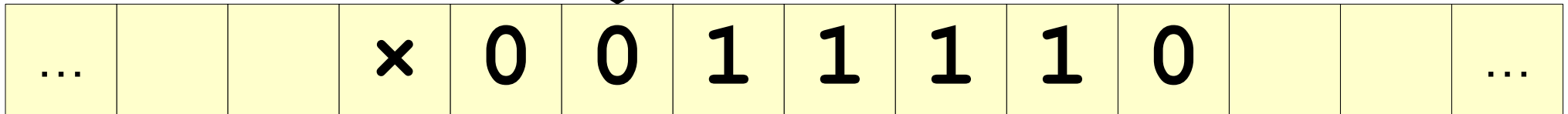
# The Solution



# The Solution



# The Solution

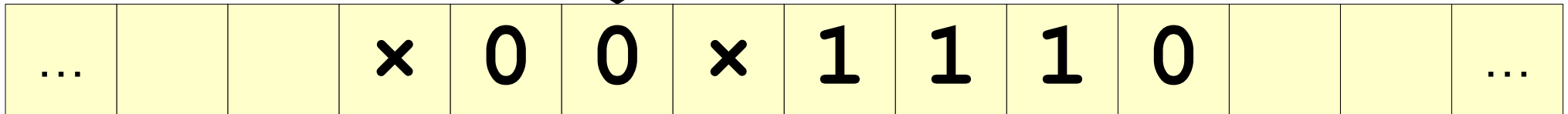


# The Solution

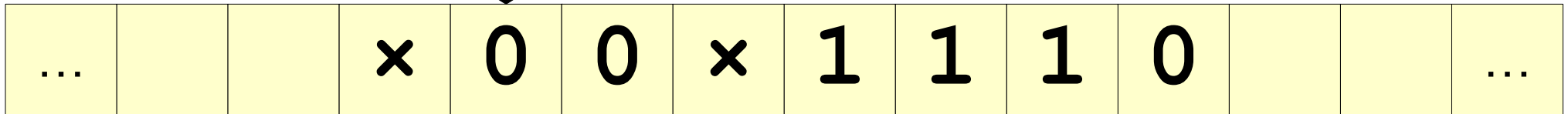


...			×	0	0	1	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

# The Solution

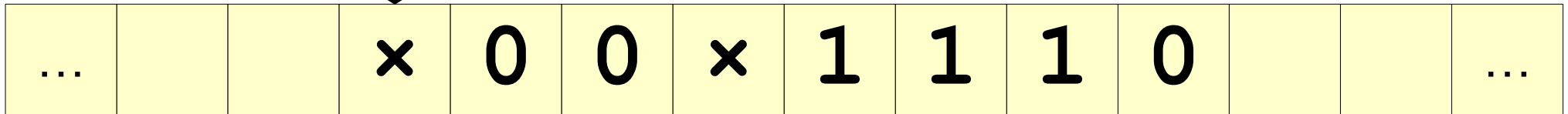


# The Solution

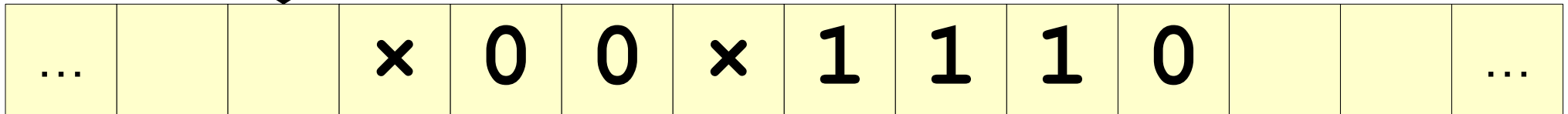




# The Solution



# The Solution

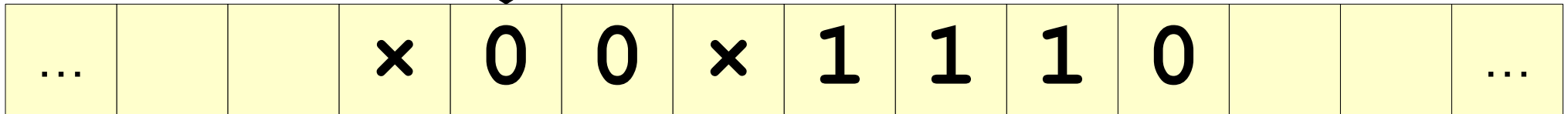


# The Solution

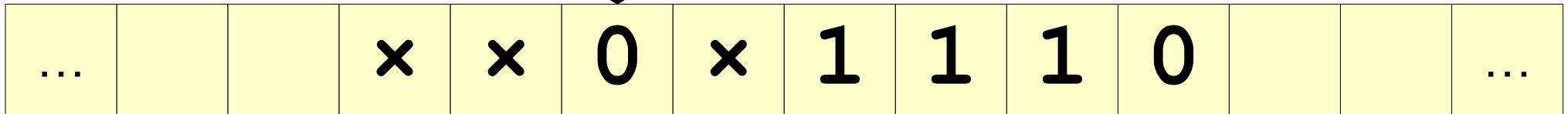


...			x	0	0	x	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

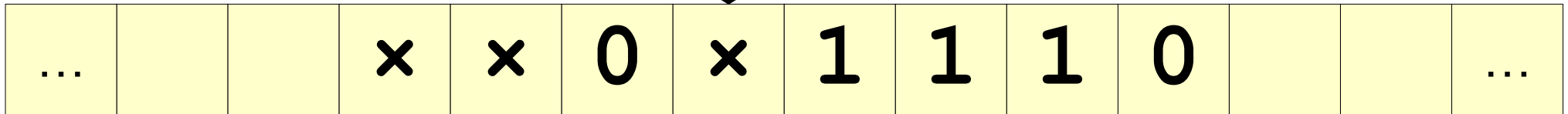
# The Solution



# The Solution



# The Solution

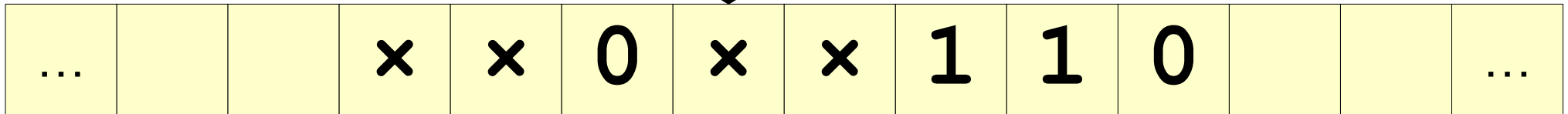


# The Solution



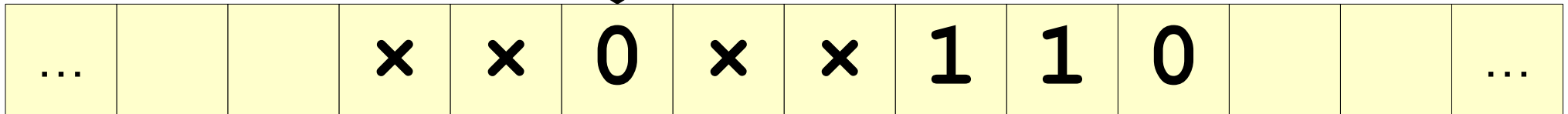
...			x	x	0	x	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

# The Solution

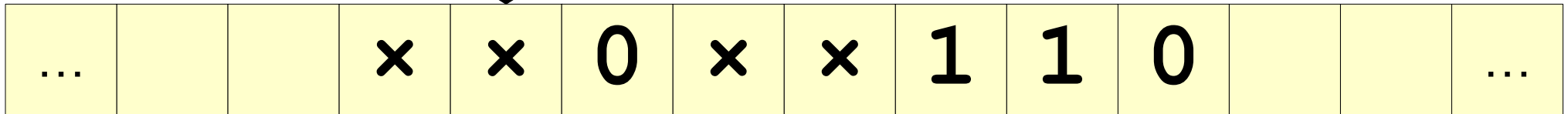




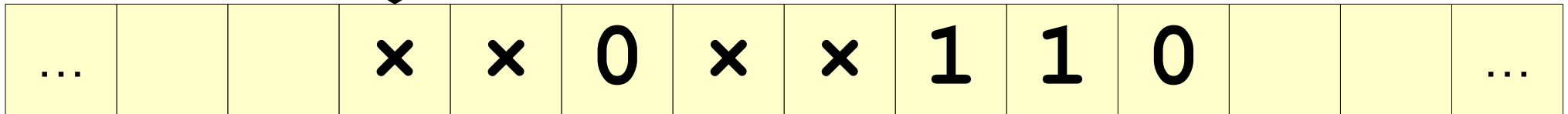
# The Solution



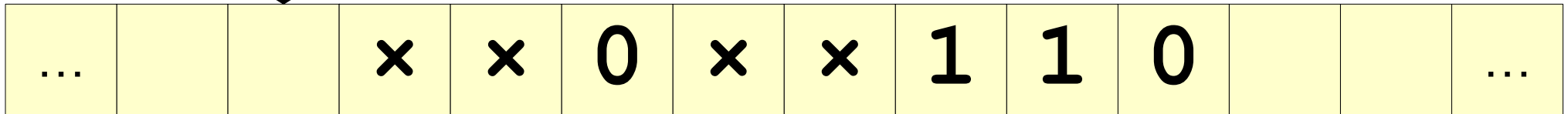
# The Solution



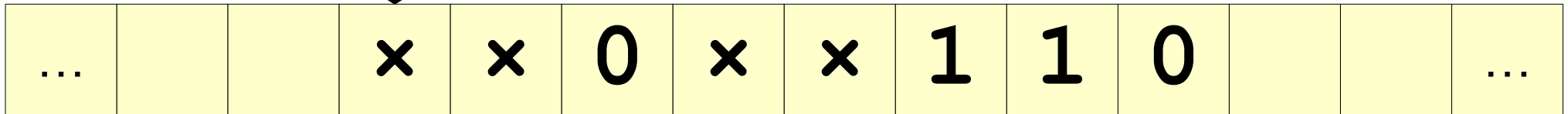
# The Solution



# The Solution

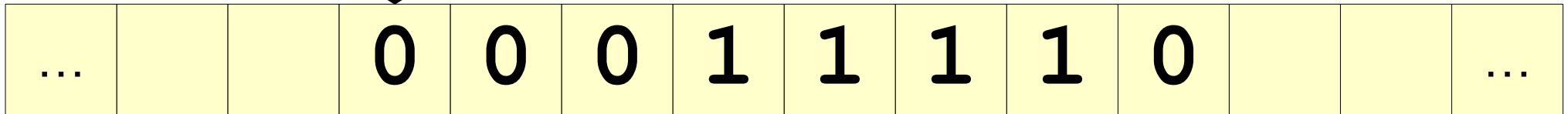


# The Solution

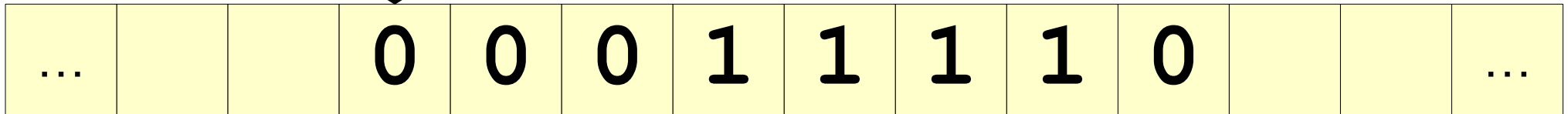


A Different Idea

# A Different Strategy



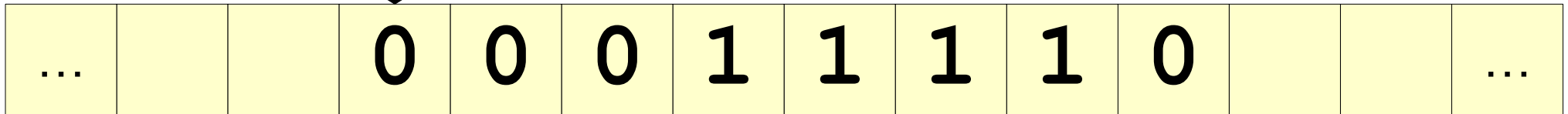
# A Different Strategy



Could we sort  
the characters of  
this string?

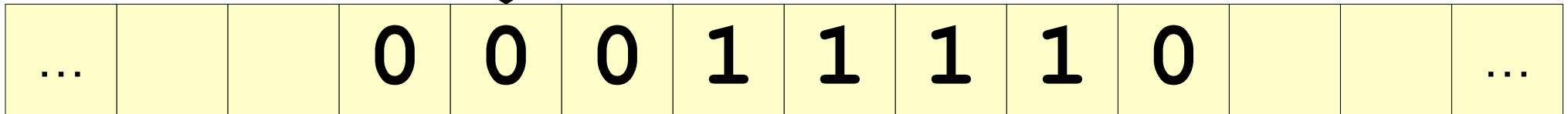


# A Different Strategy



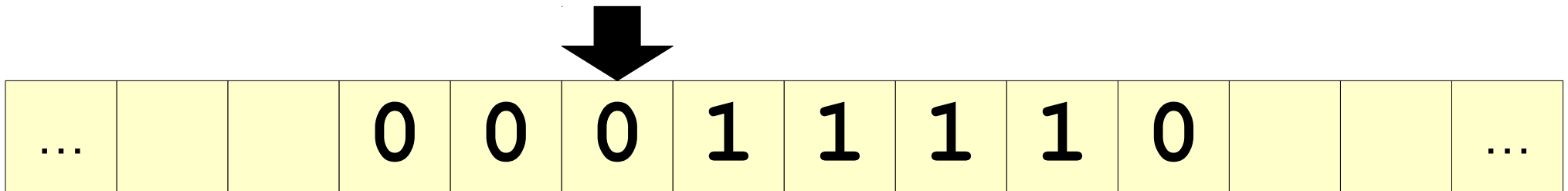
Observation 1: A string of 0s and 1s is sorted if it matches the regex  $0^*1^*$ .

# A Different Strategy



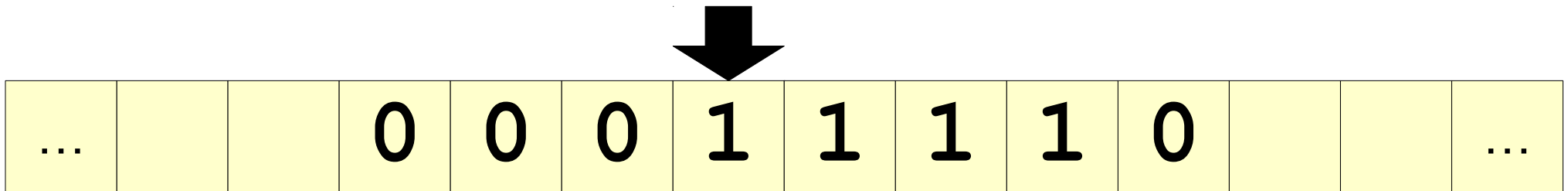
Observation 1: A string of 0s and 1s is sorted if it matches the regex  $0^*1^*$ .

# A Different Strategy



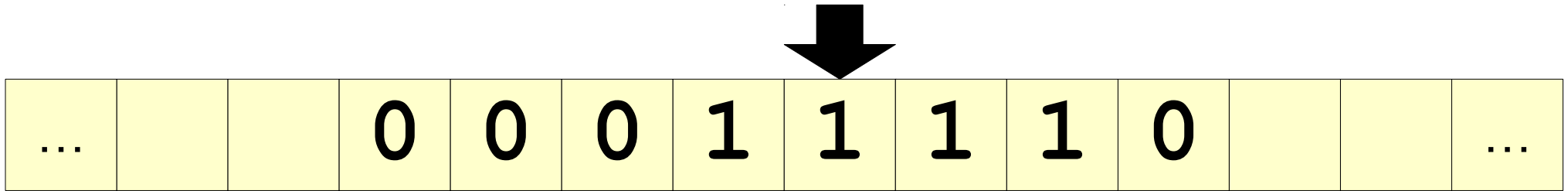
Observation 1: A string  
of 0s and 1s is sorted  
if it matches the regex  
 $0^*1^*$ .

# A Different Strategy



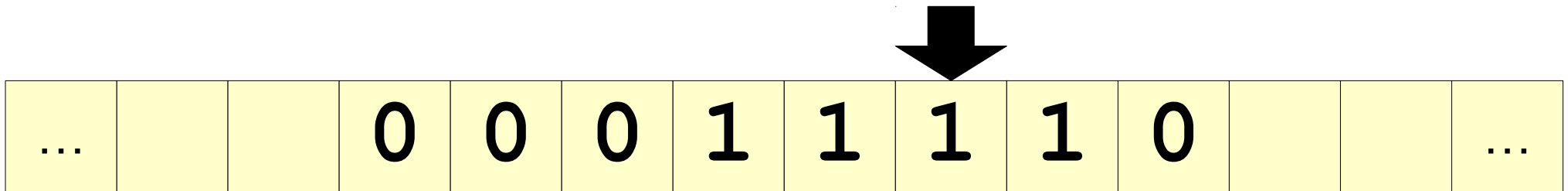
Observation 1: A string  
of 0s and 1s is sorted  
if it matches the regex  
 $0^*1^*$ .

# A Different Strategy



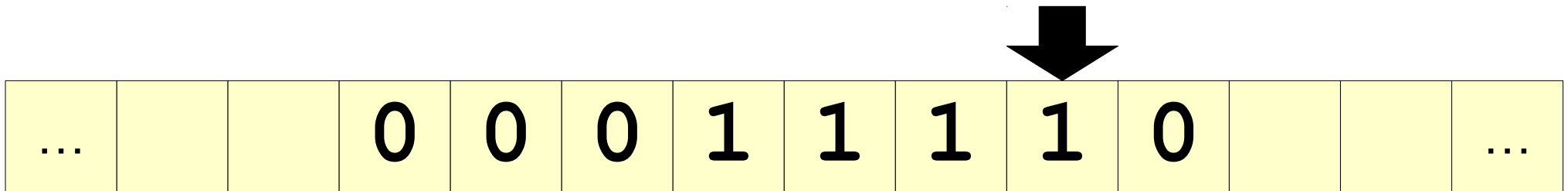
Observation 1: A string  
of 0s and 1s is sorted  
if it matches the regex  
 $0^*1^*$ .

# A Different Strategy



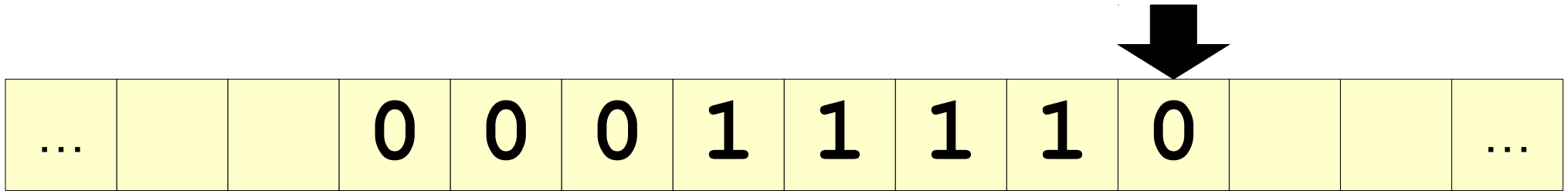
Observation 1: A string  
of 0s and 1s is sorted  
if it matches the regex  
 $0^*1^*$ .

# A Different Strategy



Observation 1: A string  
of 0s and 1s is sorted  
if it matches the regex  
 $0^*1^*$ .

# A Different Strategy



Observation 1: A string  
of 0s and 1s is sorted  
if it matches the regex  
 $0^*1^*$ .



# A Different Strategy



Observation 1: A string  
of 0s and 1s is sorted  
if it matches the regex  
 $0^*1^*$ .

# A Different Strategy



Observation 2: A string of 0s and 1s is not sorted if it contains 10 as a substring.

# A Different Strategy



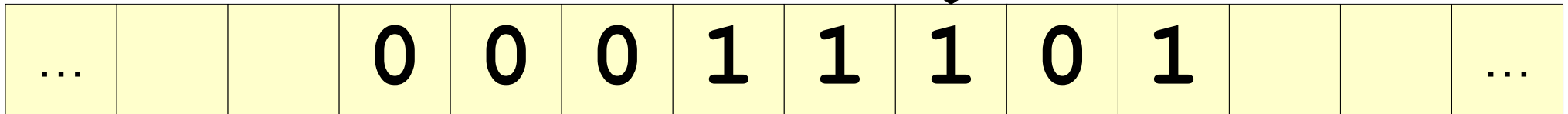
Observation 2: A string of 0s and 1s is not sorted if it contains 10 as a substring.

# A Different Strategy



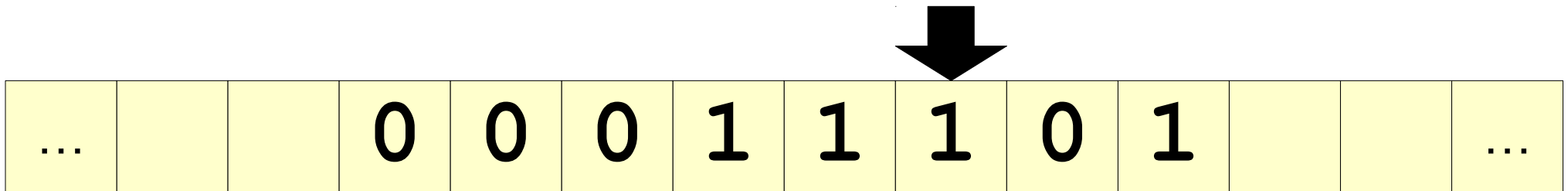
Observation 2: A string of 0s and 1s is not sorted if it contains 10 as a substring.

# A Different Strategy



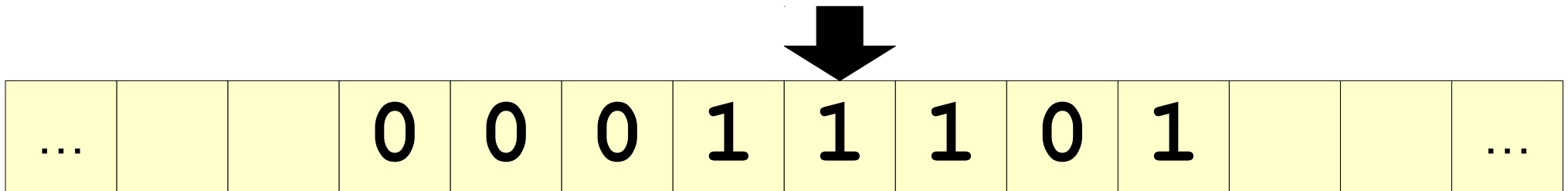
Observation 2: A string of 0s and 1s is not sorted if it contains 10 as a substring.

# A Different Strategy



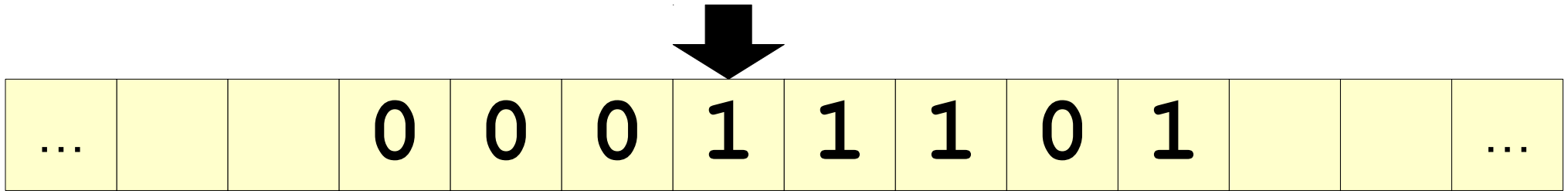
**Idea:** Repeatedly find a copy of 10 and replace it with 01.

# A Different Strategy



**Idea:** Repeatedly find a copy of 10 and replace it with 01.

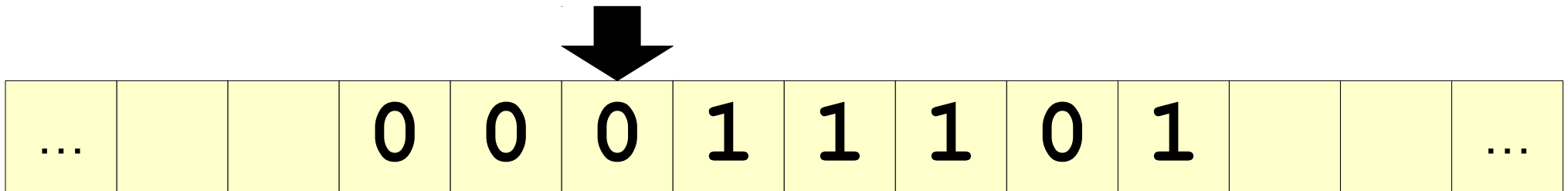
# A Different Strategy



**Idea:** Repeatedly find a copy of 10 and replace it with 01.

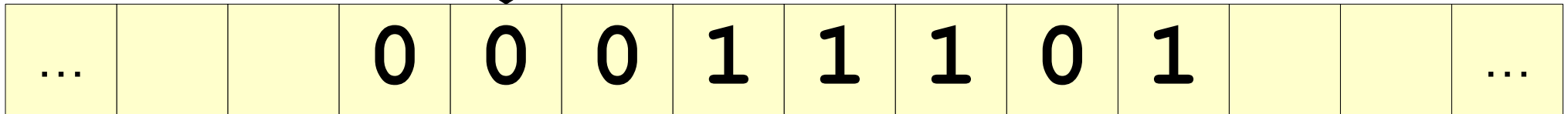


# A Different Strategy



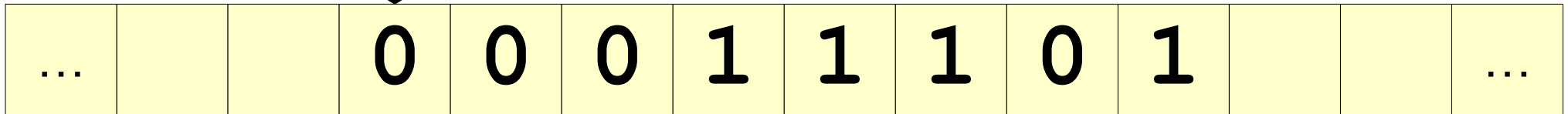
**Idea:** Repeatedly find a copy of 10 and replace it with 01.

# A Different Strategy



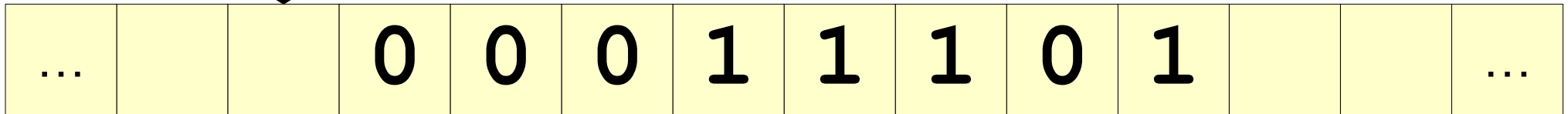
**Idea:** Repeatedly find a copy of 10 and replace it with 01.

# A Different Strategy



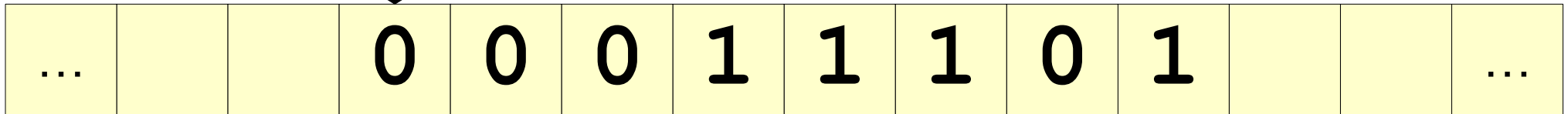
**Idea:** Repeatedly find a copy of 10 and replace it with 01.

# A Different Strategy



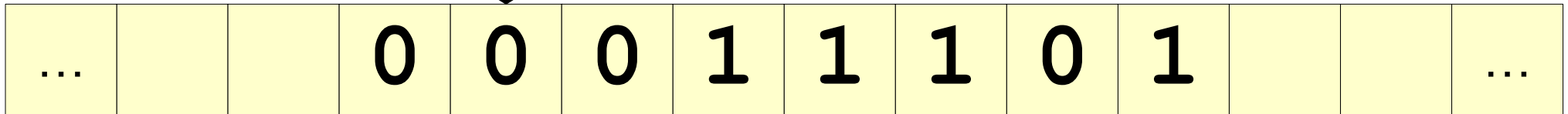
**Idea:** Repeatedly find a copy of 10 and replace it with 01.

# A Different Strategy



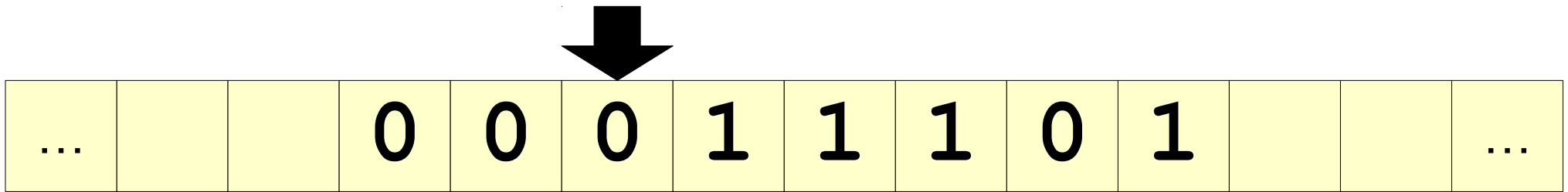
**Idea:** Repeatedly find a copy of 10 and replace it with 01.

# A Different Strategy



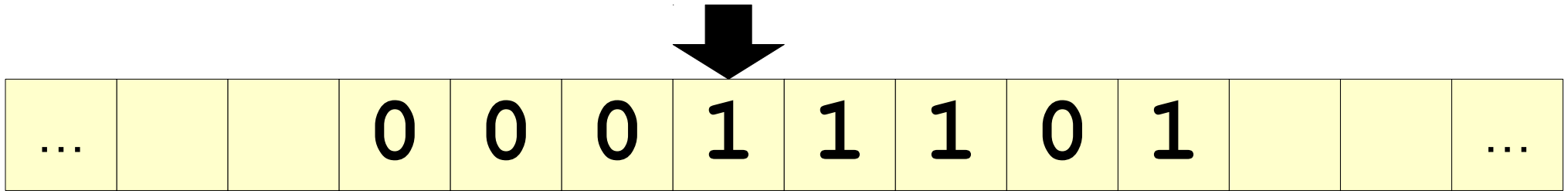
**Idea:** Repeatedly find a copy of 10 and replace it with 01.

# A Different Strategy



**Idea:** Repeatedly find a copy of 10 and replace it with 01.

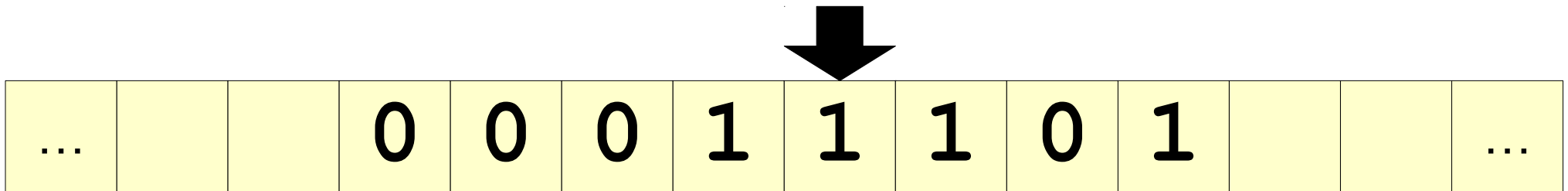
# A Different Strategy



**Idea:** Repeatedly find a copy of 10 and replace it with 01.

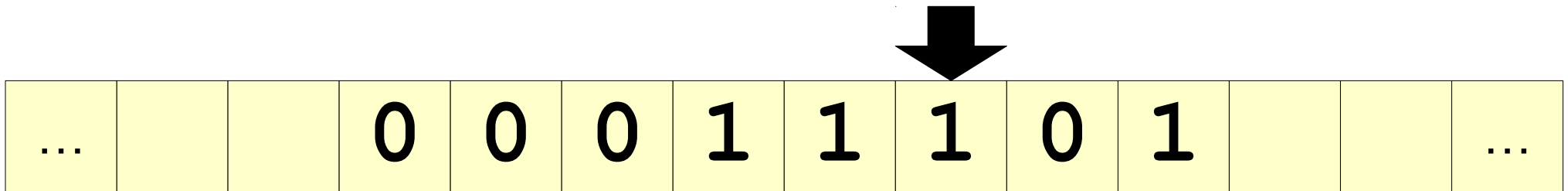


# A Different Strategy



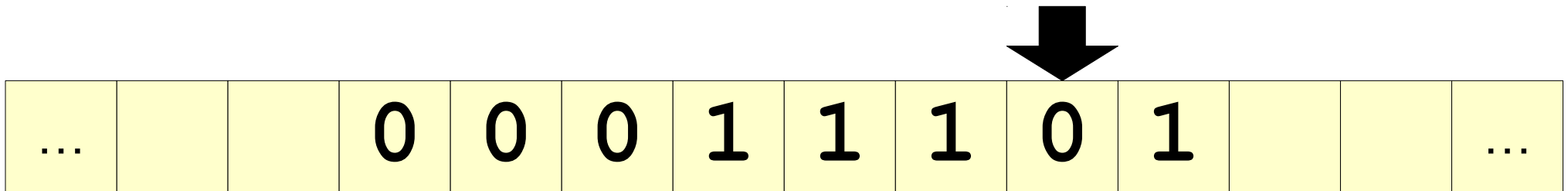
**Idea:** Repeatedly find a copy of 10 and replace it with 01.

# A Different Strategy



**Idea:** Repeatedly find a copy of 10 and replace it with 01.

# A Different Strategy



**Idea:** Repeatedly find a copy of 10 and replace it with 01.

# A Different Strategy



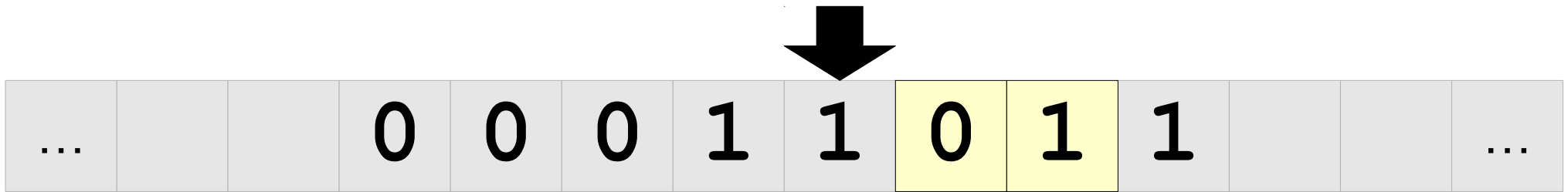
**Idea:** Repeatedly find a copy of 10 and replace it with 01.

# A Different Strategy



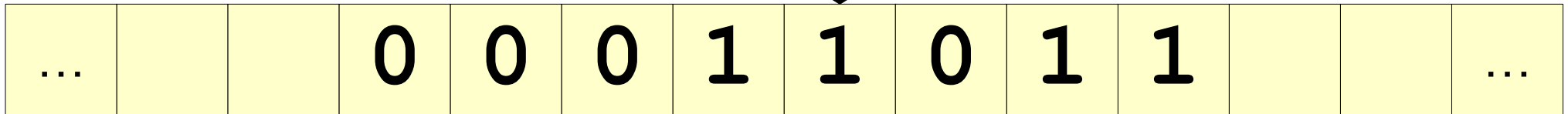
**Idea:** Repeatedly find a copy of 10 and replace it with 01.

# A Different Strategy



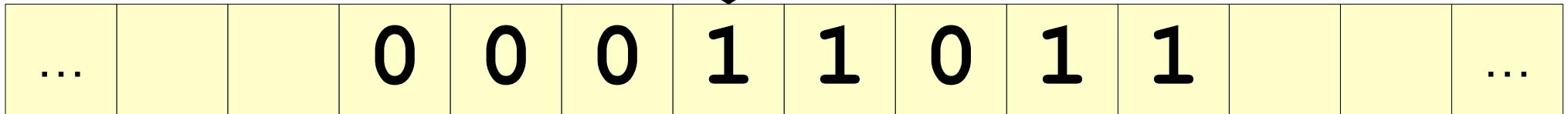
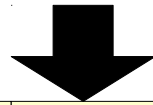
**Idea:** Repeatedly find a copy of 10 and replace it with 01.

# A Different Strategy



**Idea:** Repeatedly find a copy of 10 and replace it with 01.

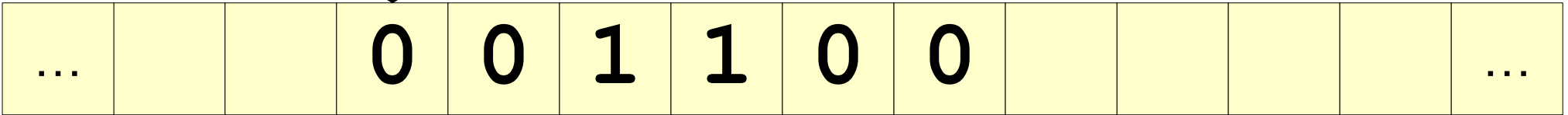
# A Different Strategy

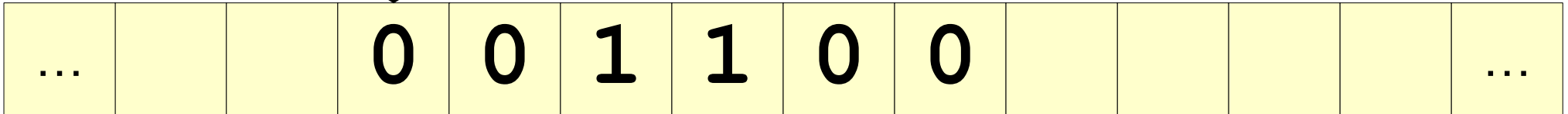
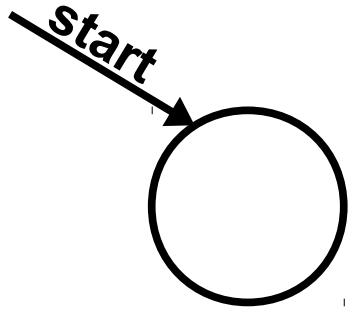


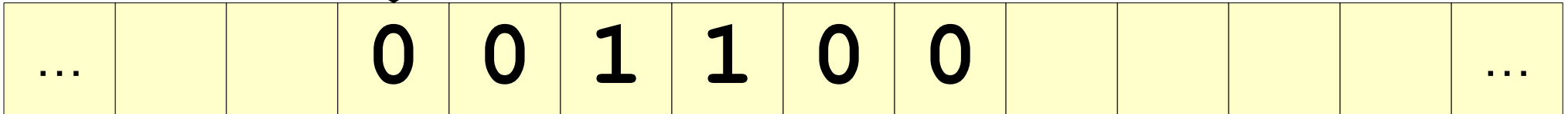
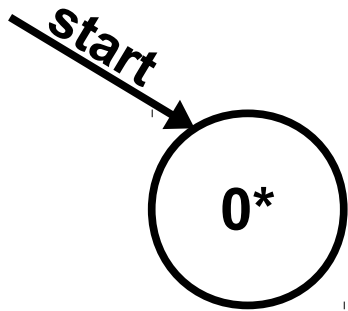
**Idea:** Repeatedly find a copy of 10 and replace it with 01.

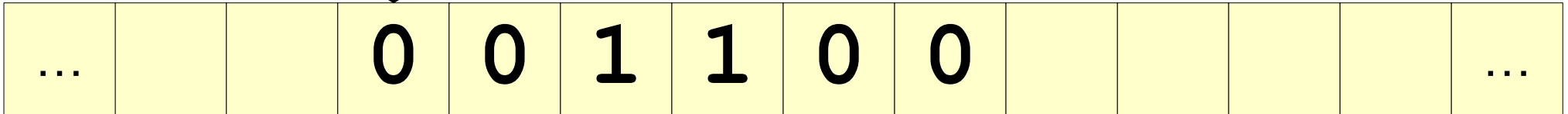
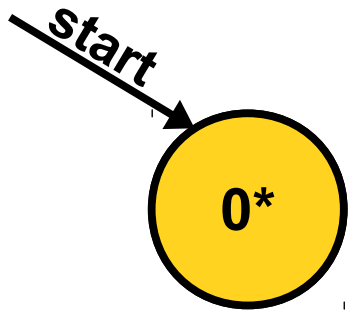


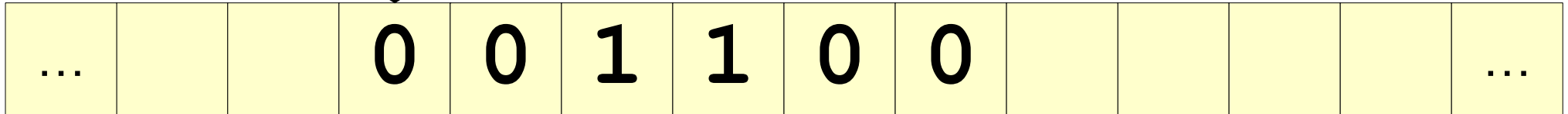
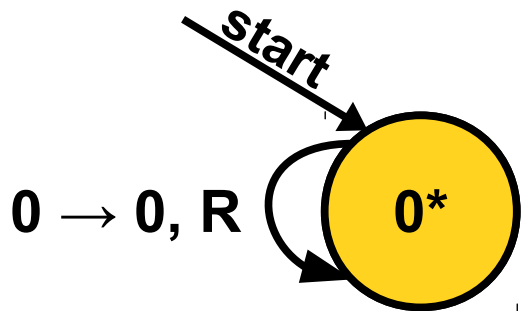
Let's Build It!

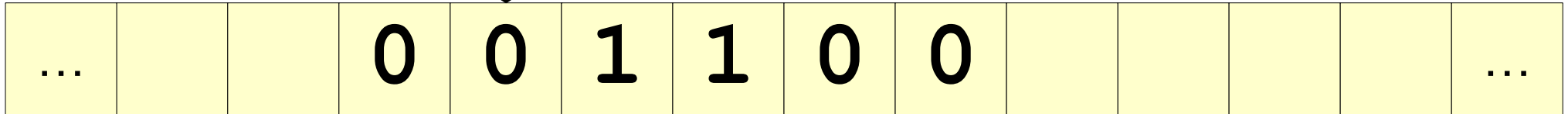
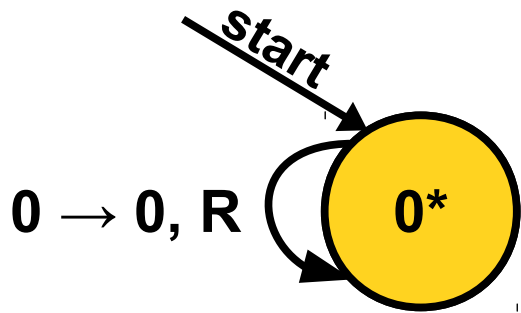


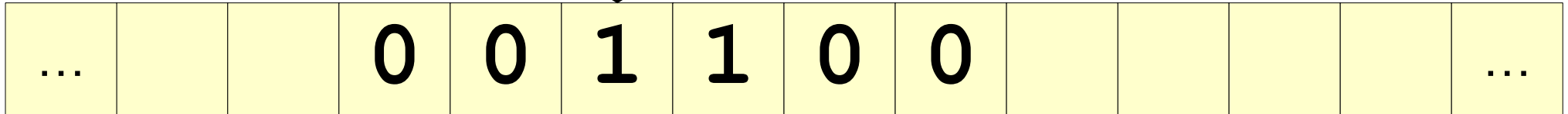
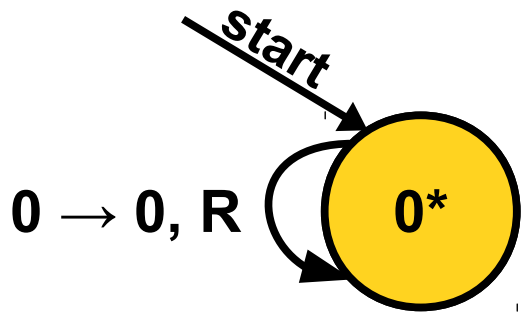




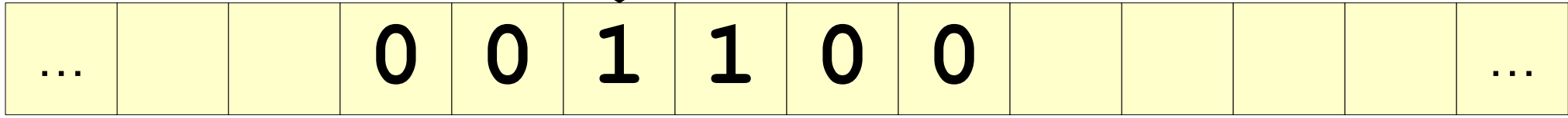
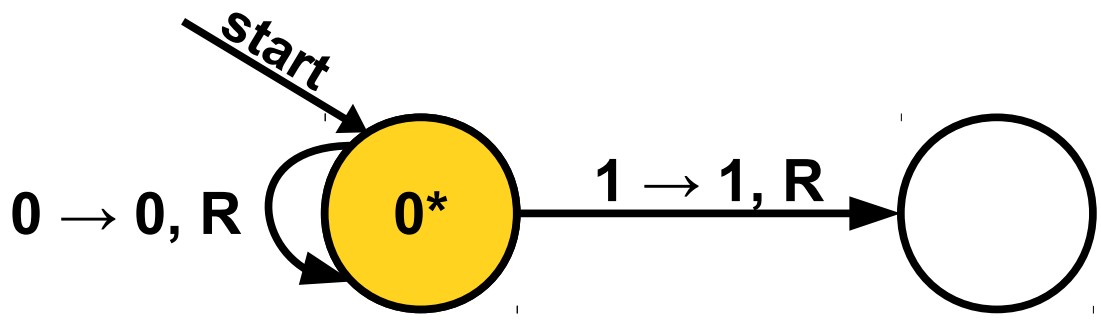


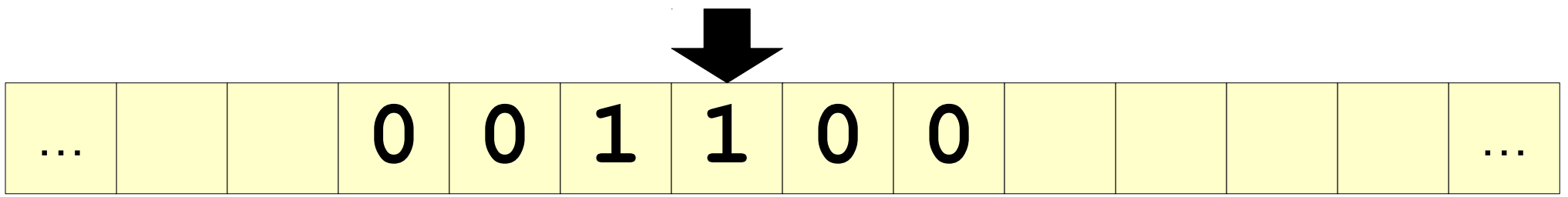
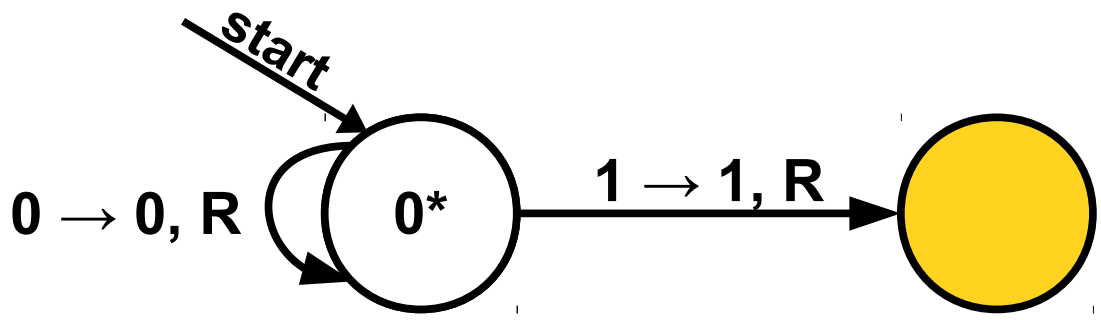


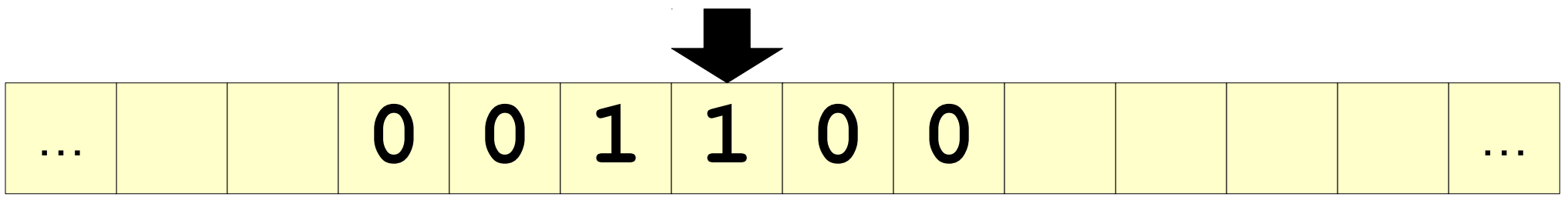
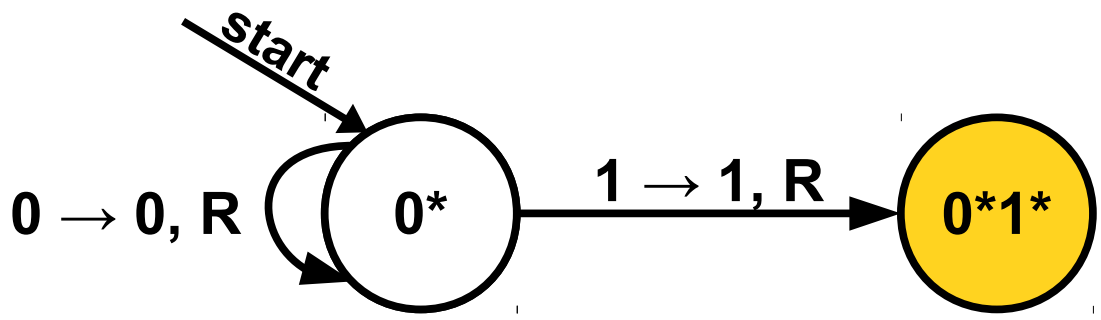


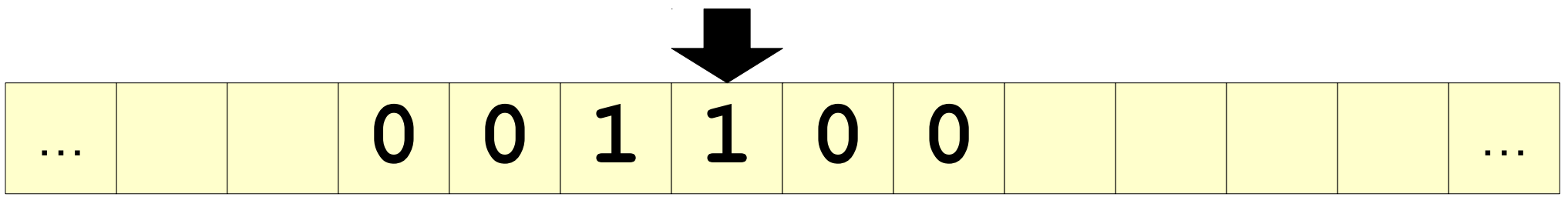
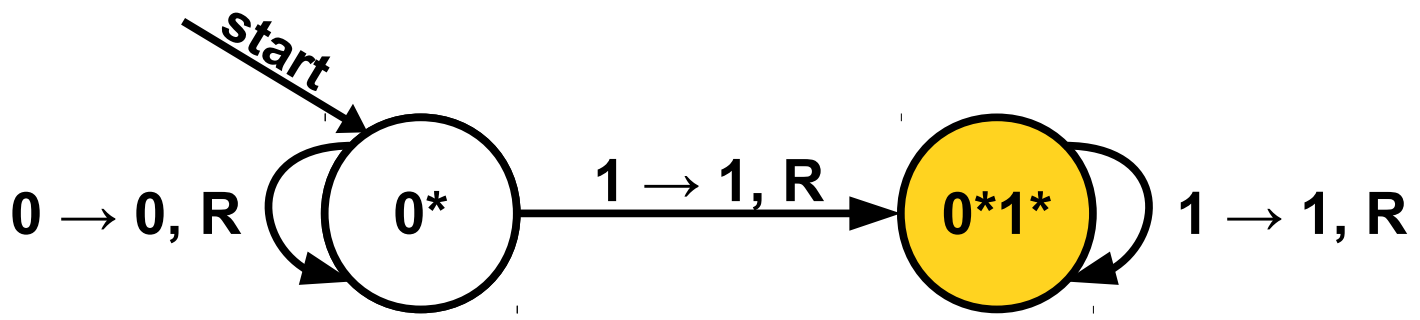


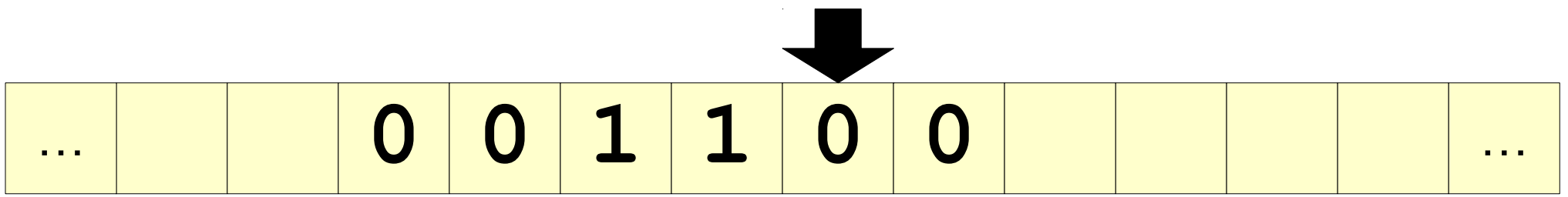
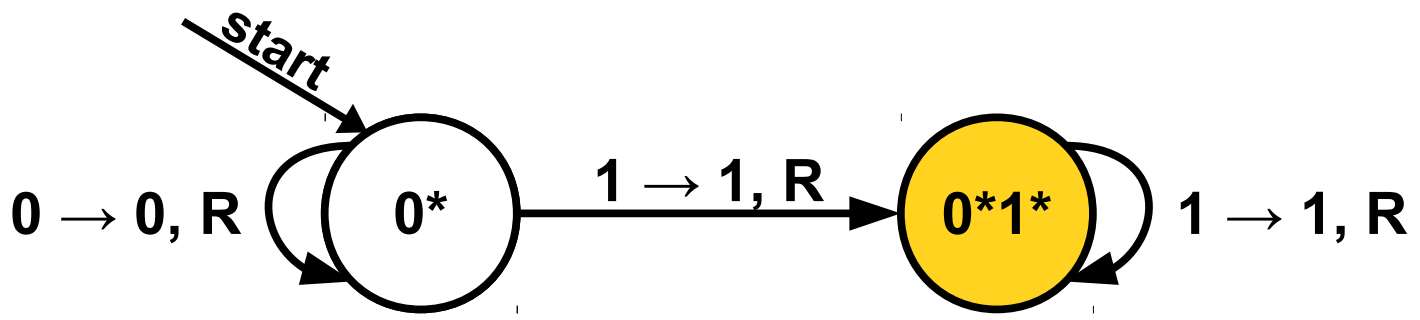


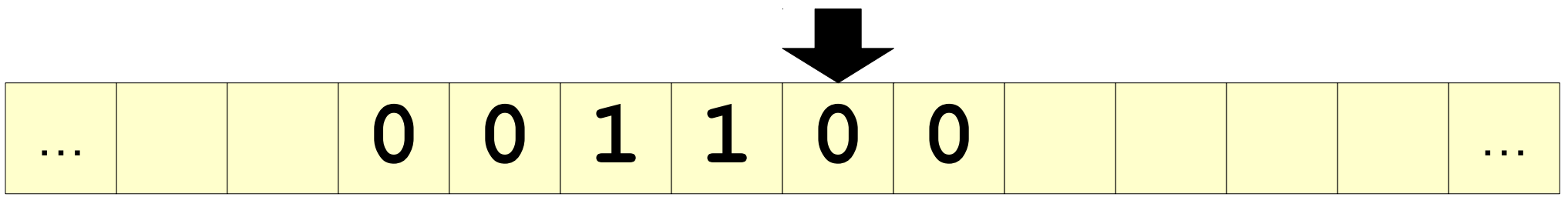
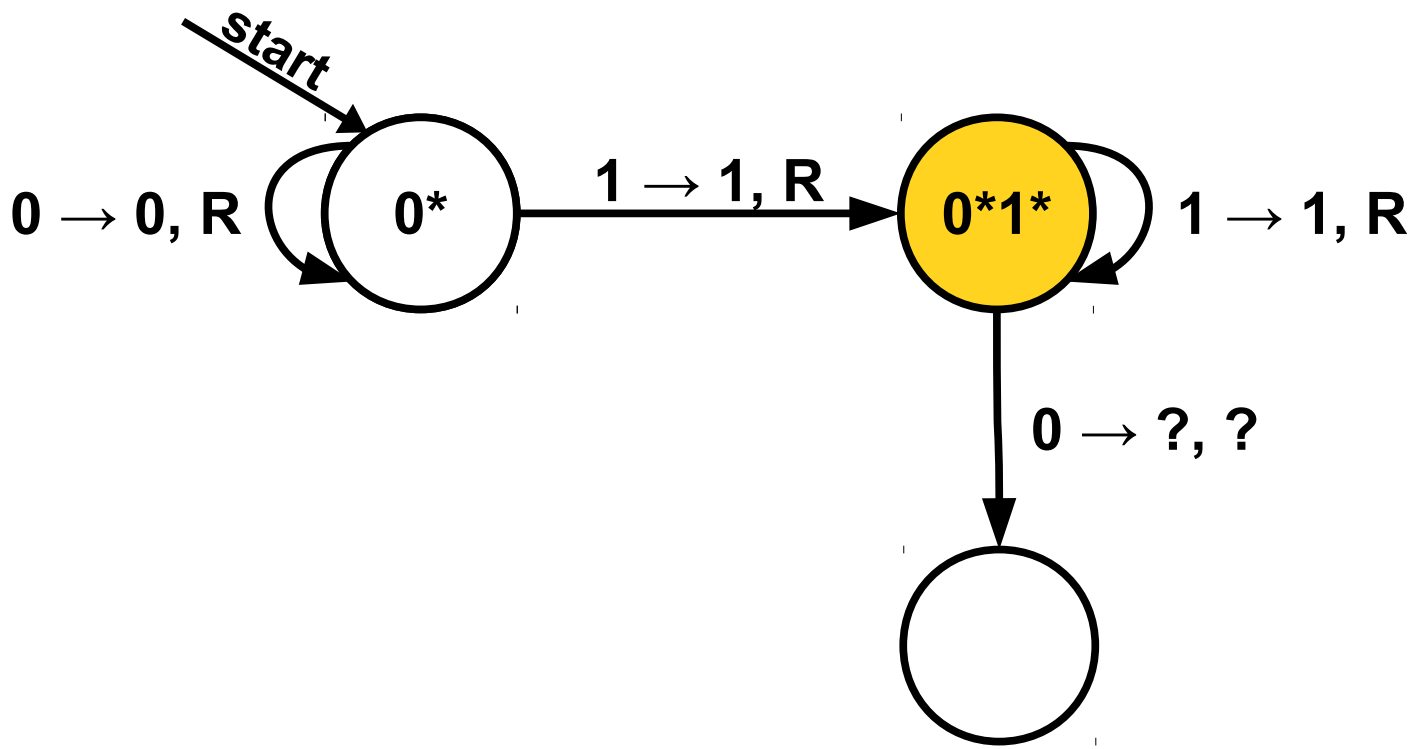






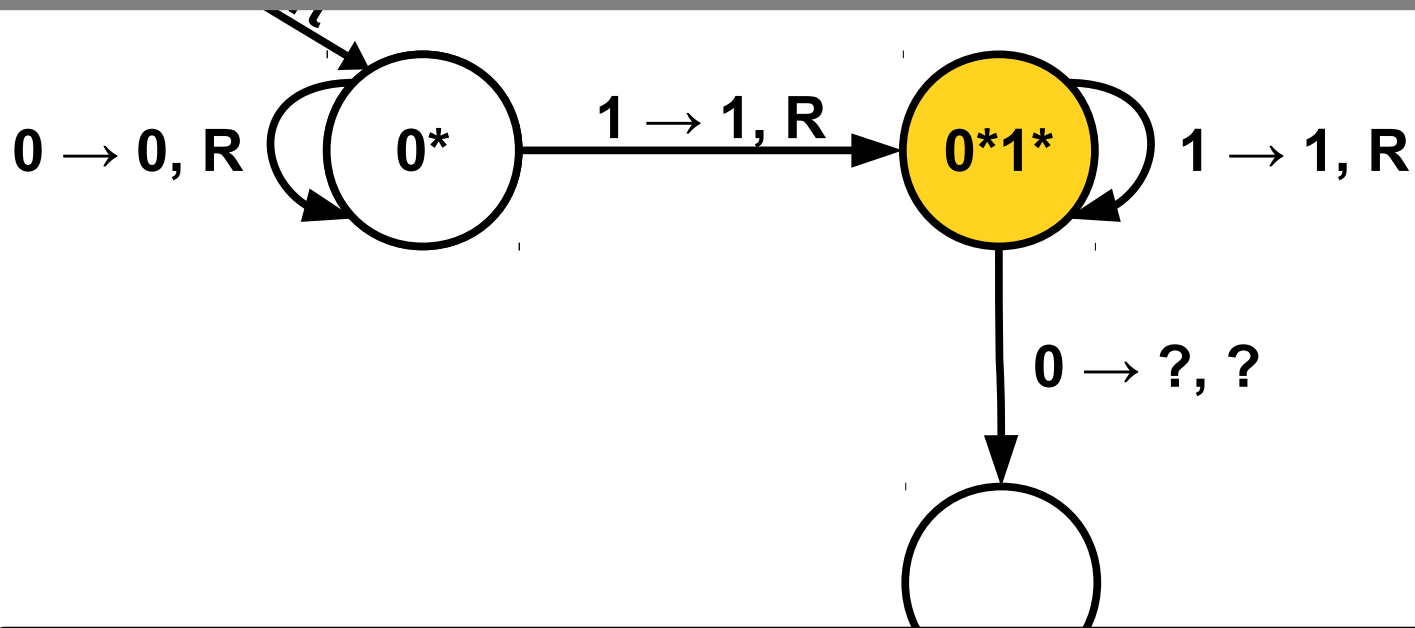




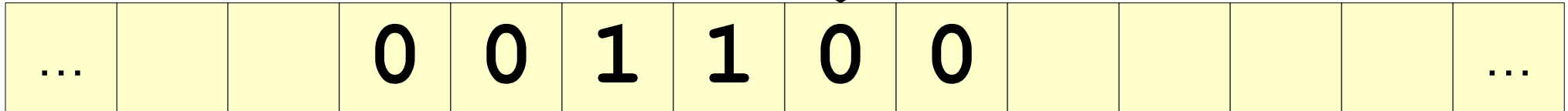


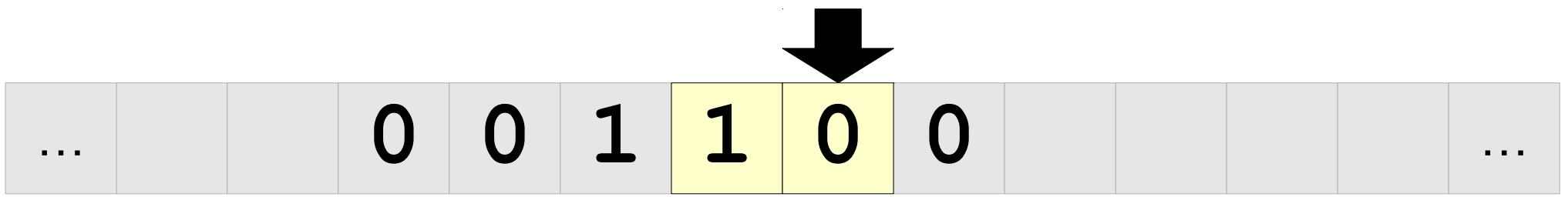
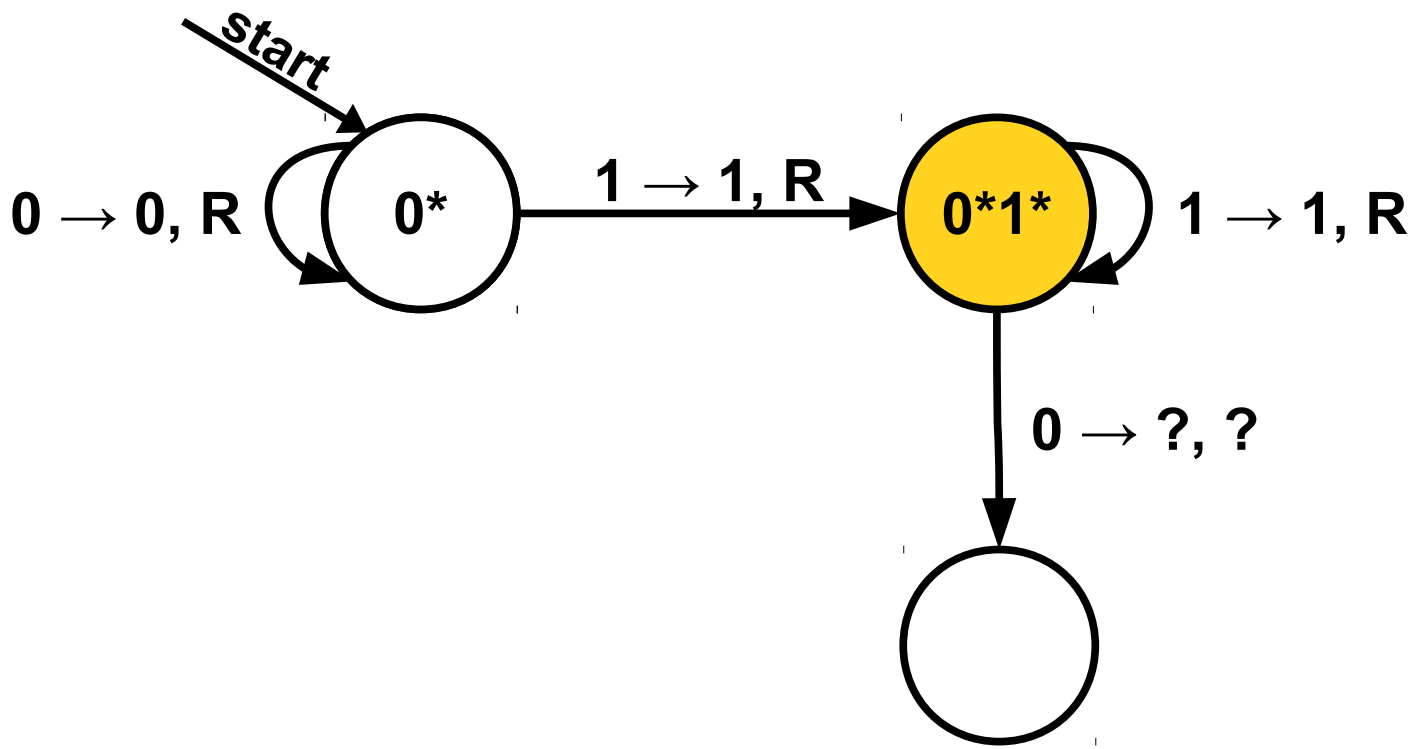
Based on we want this TM to do, what should this transition say?

- A.  $0 \rightarrow 0, R$
- B.  $0 \rightarrow 1, R$
- C.  $0 \rightarrow 0, L$
- D.  $0 \rightarrow 1, L$
- E. None of these, or two or more of these.

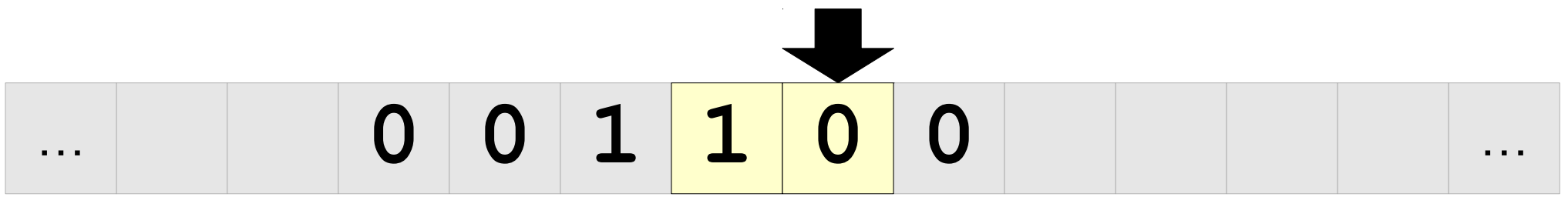
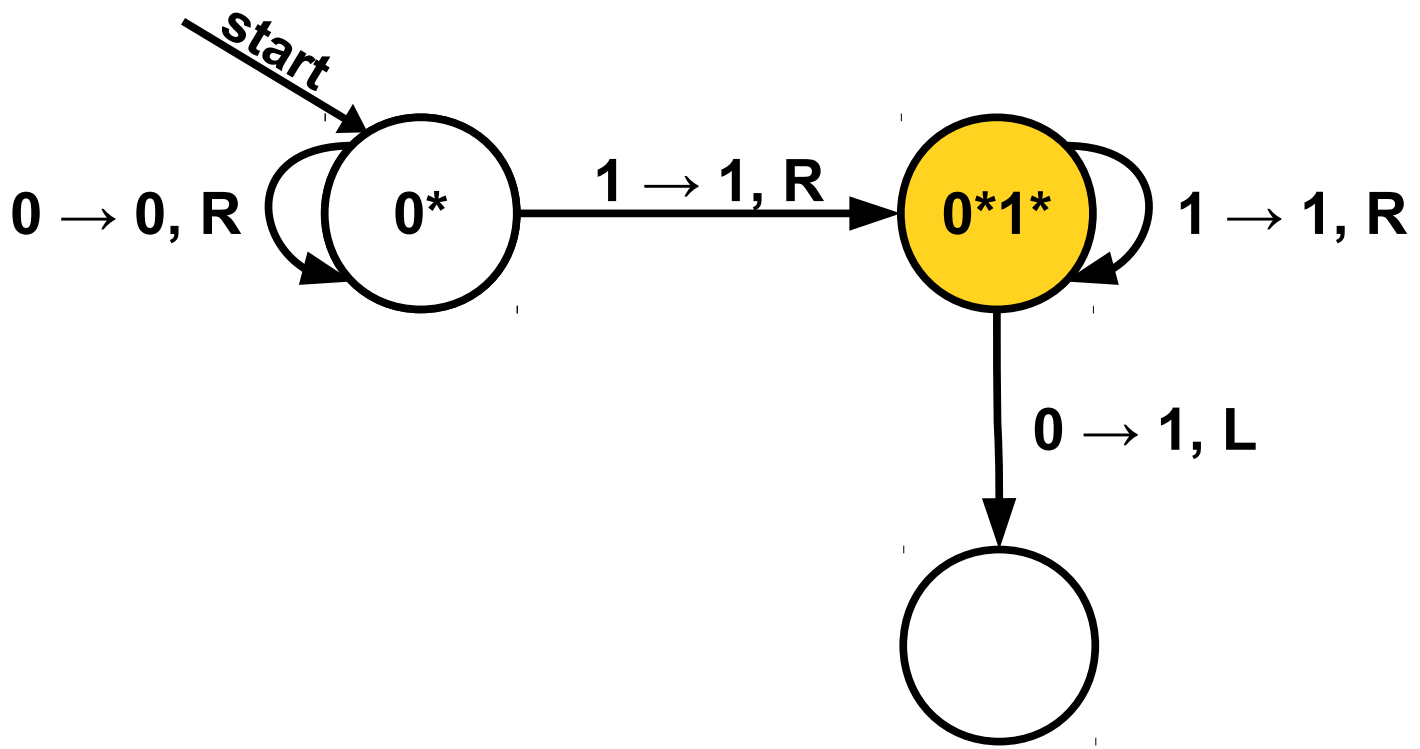


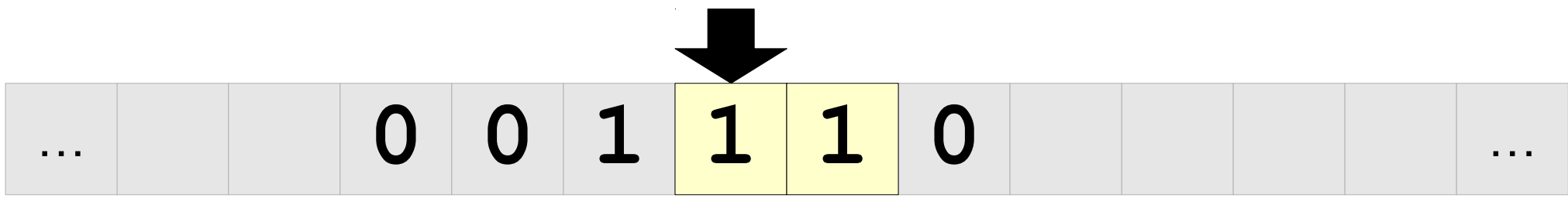
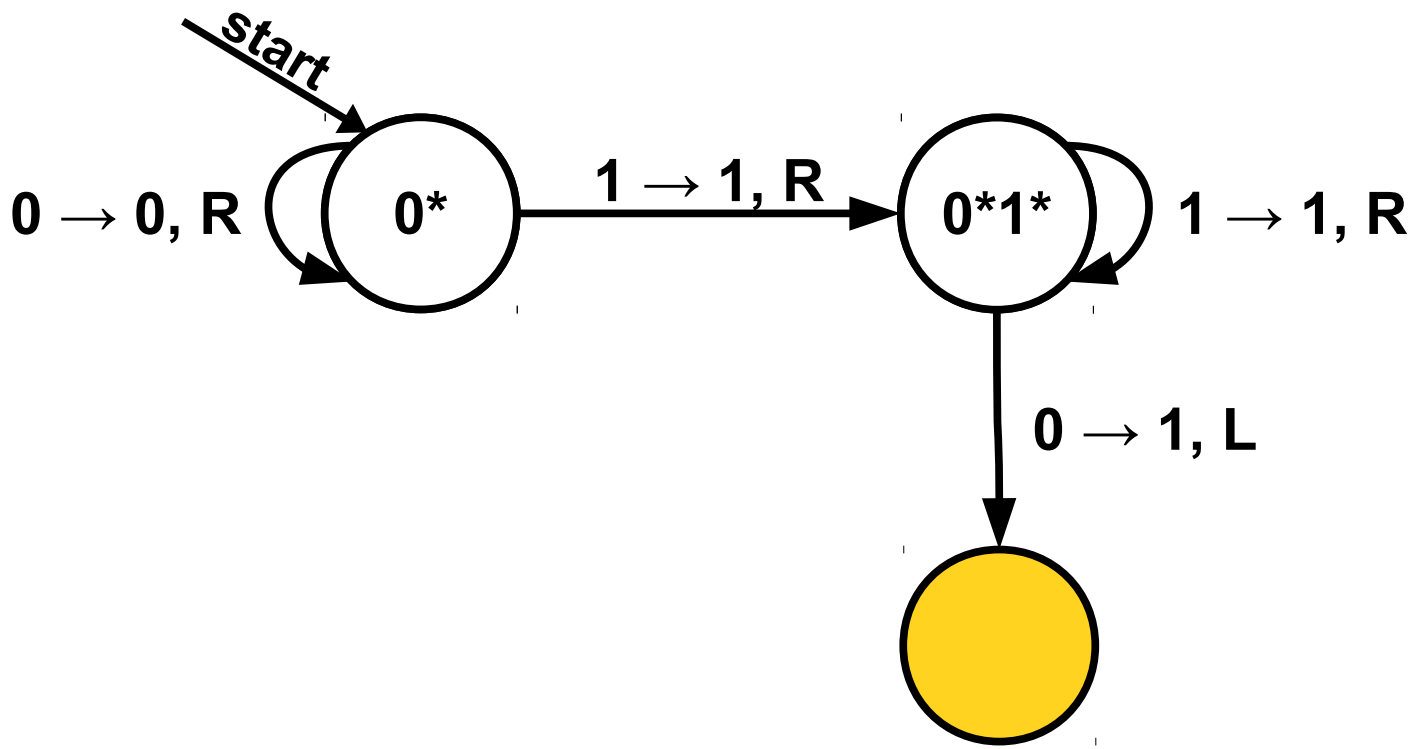
Answer at [Pollev.com/cs103](https://pollev.com/cs103) or  
text **CS103** to **22333** once to join, then **A, B, C, D, or E.**

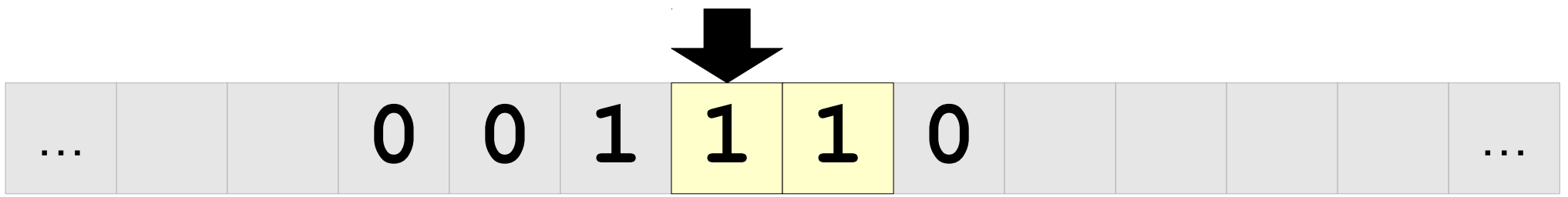
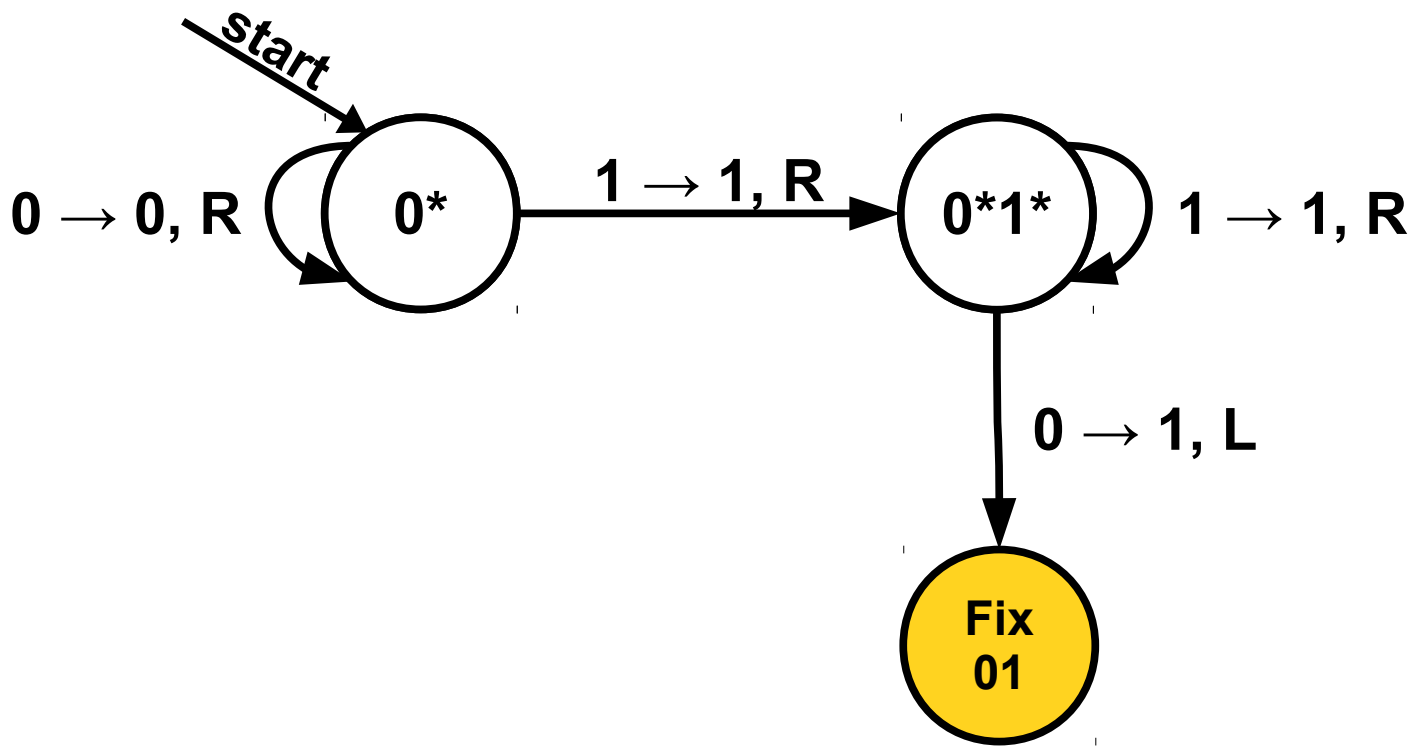


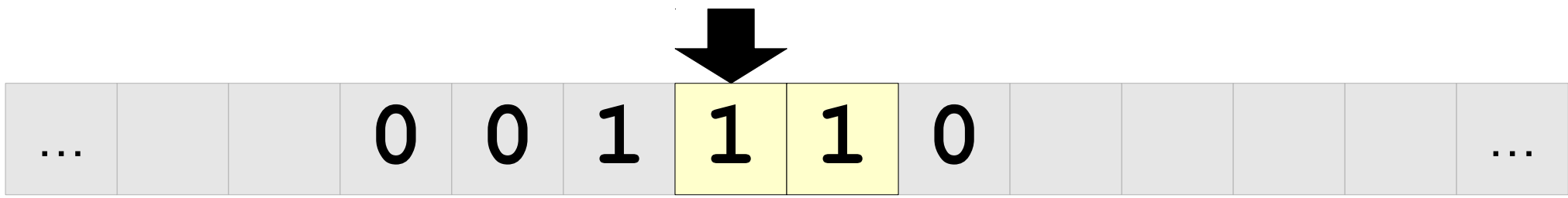
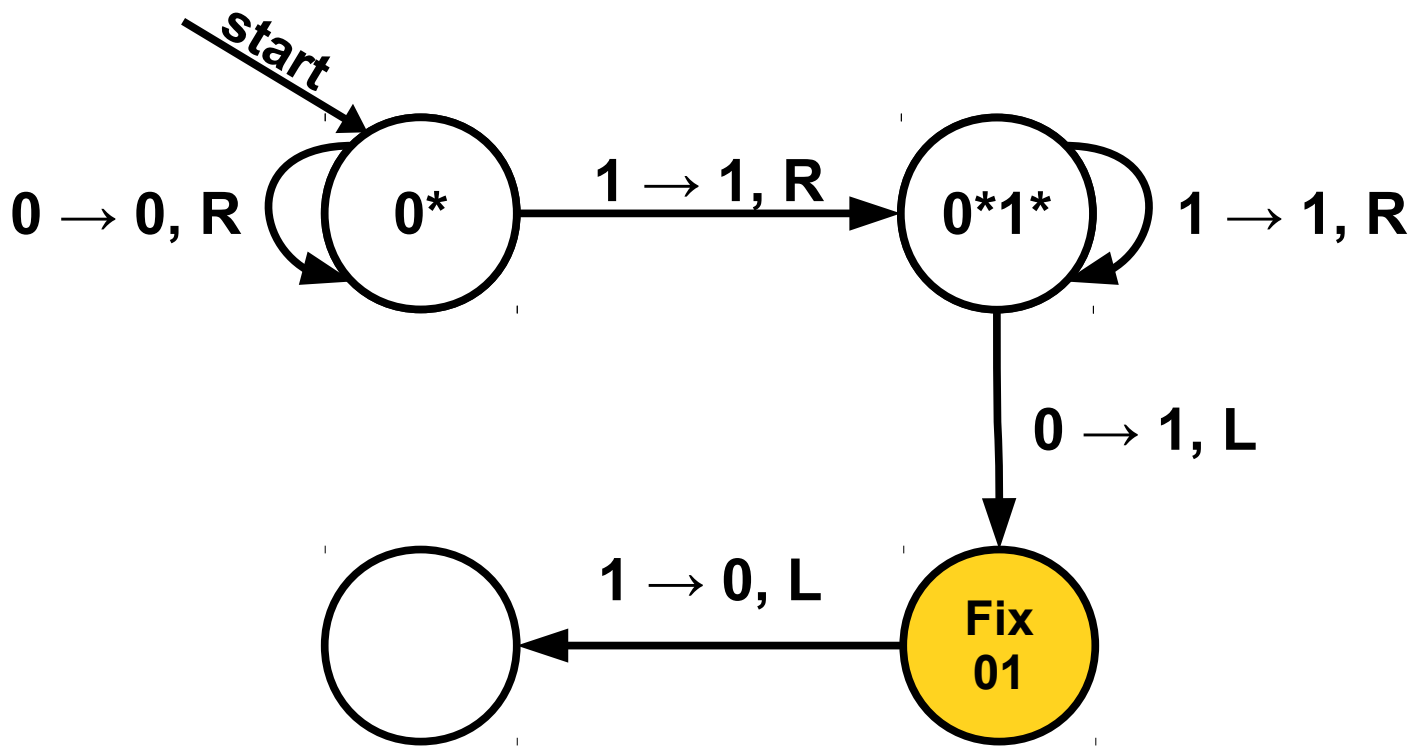




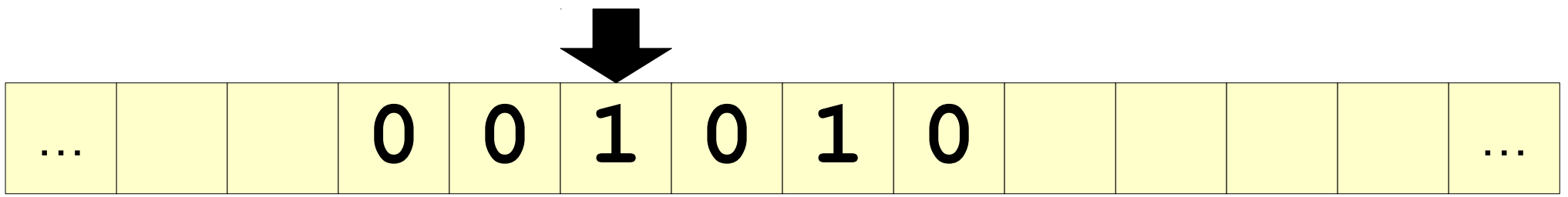
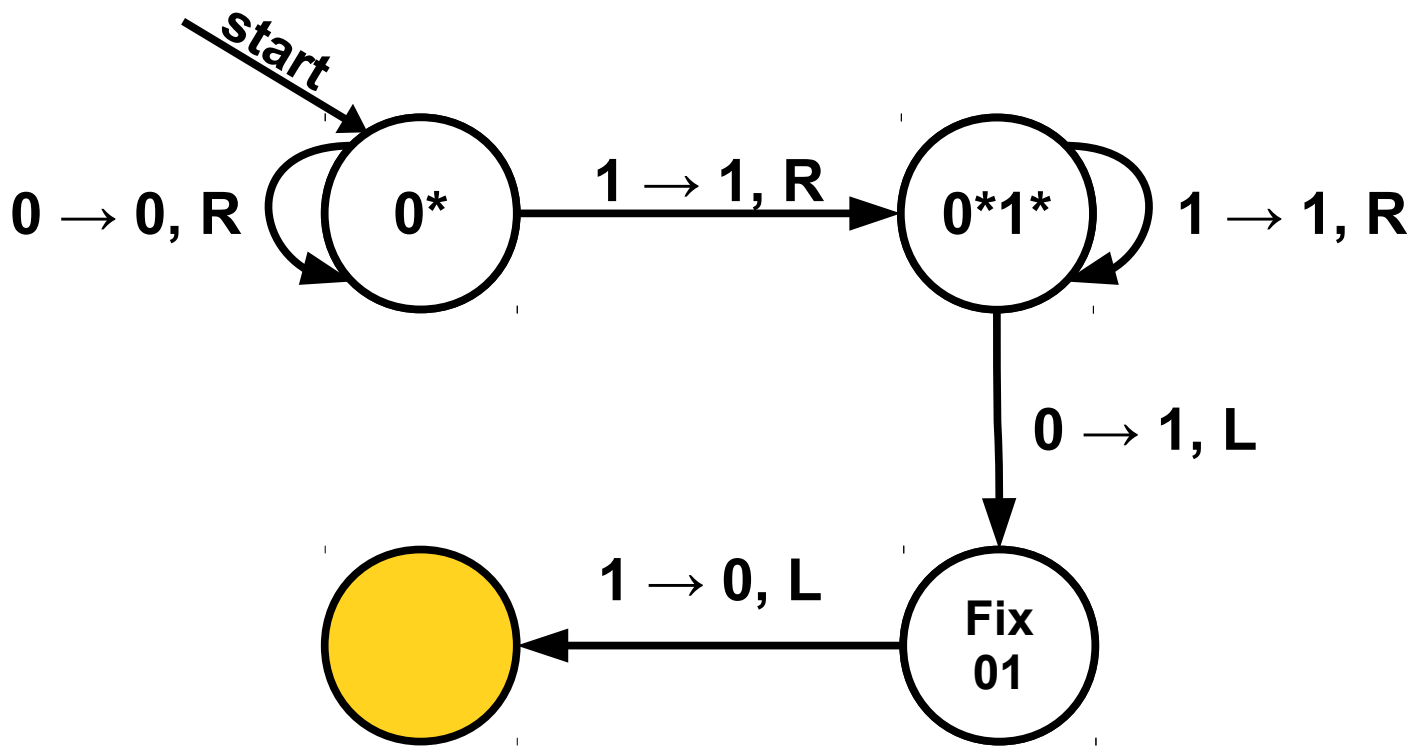


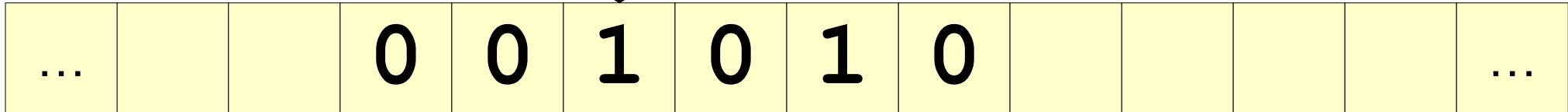
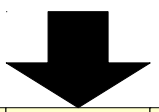
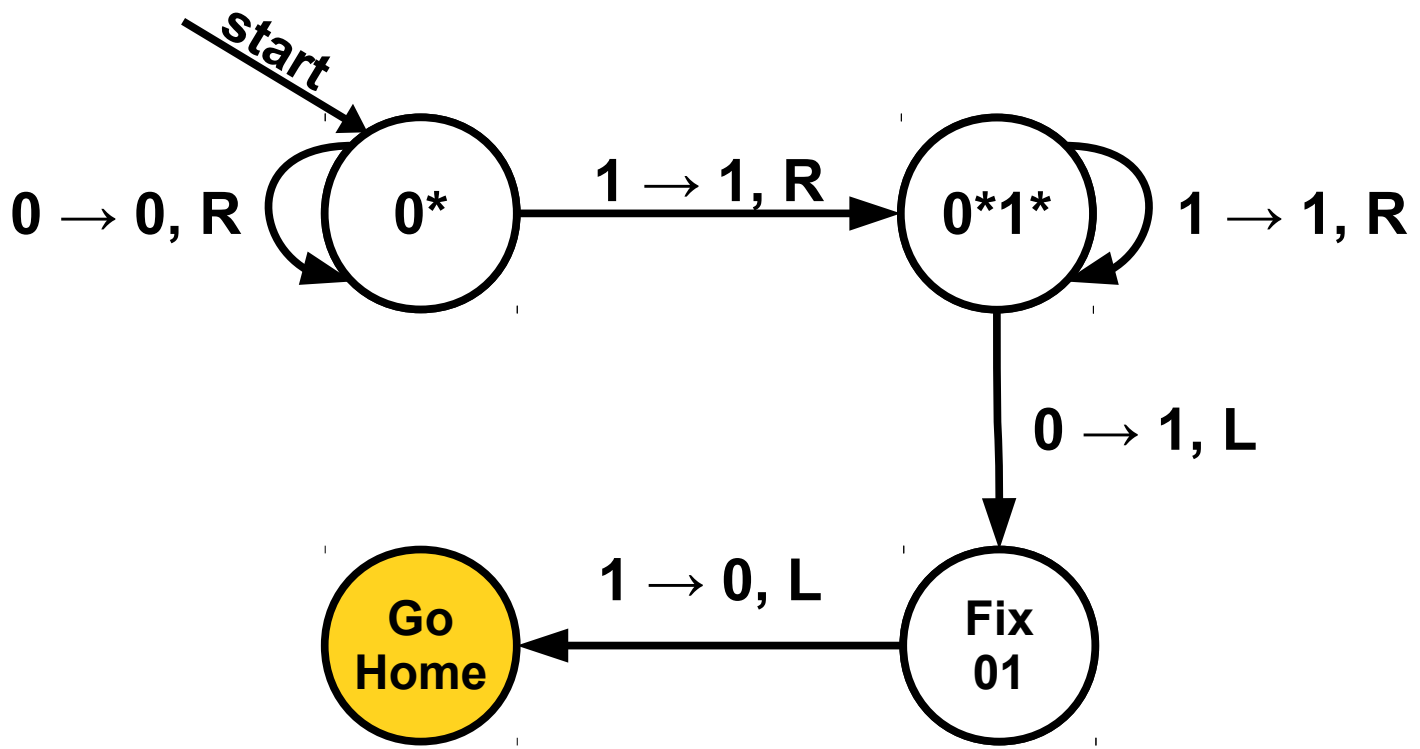


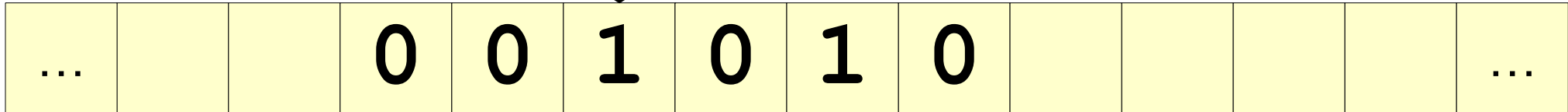
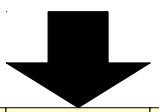
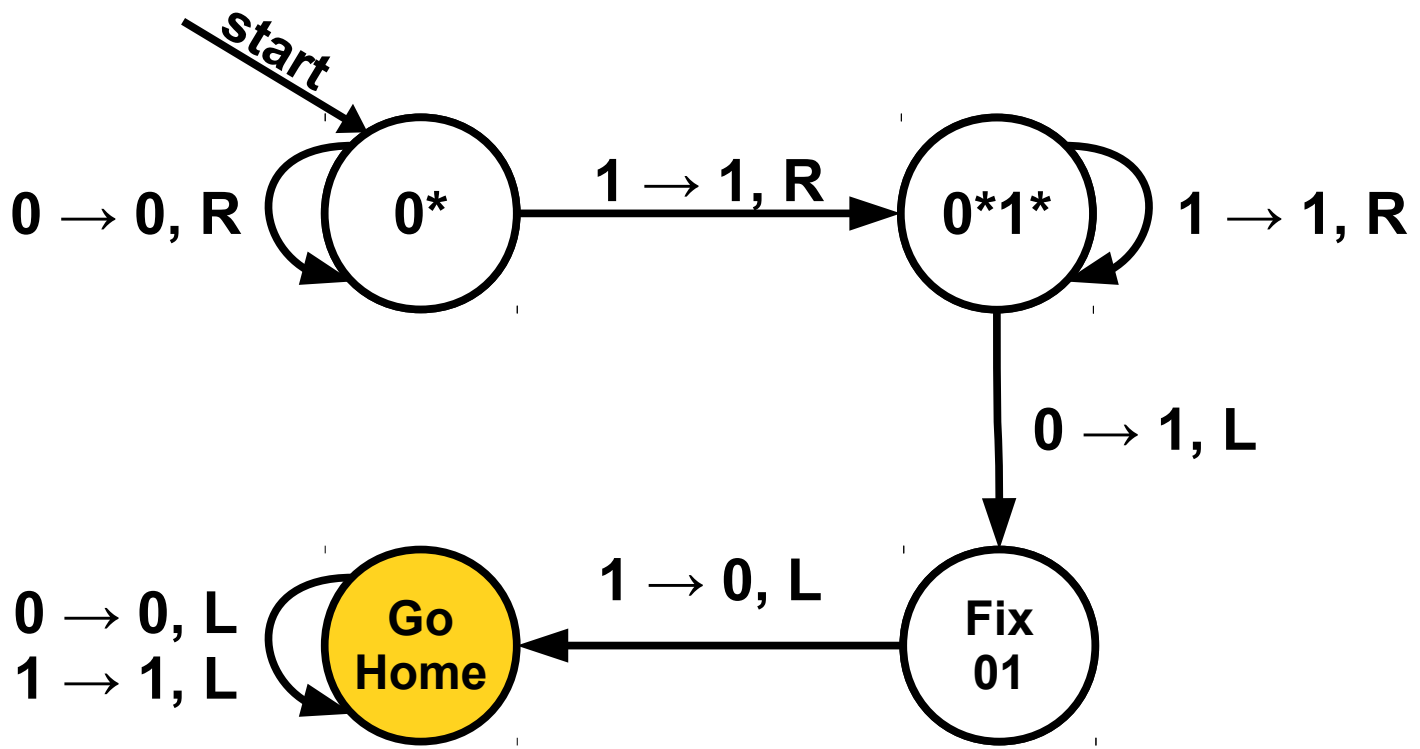




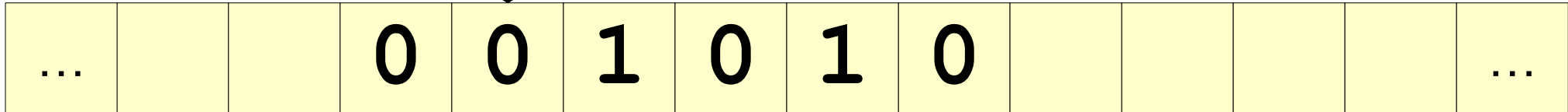
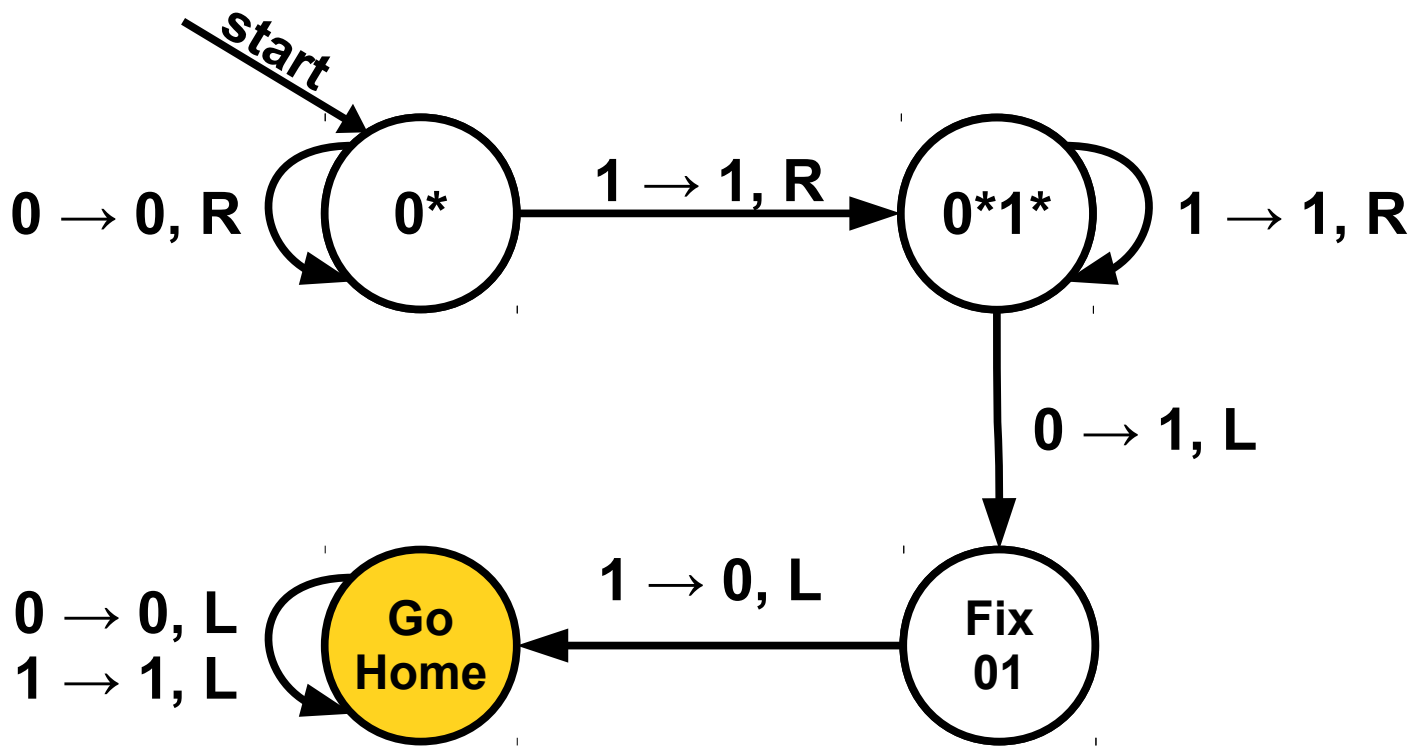


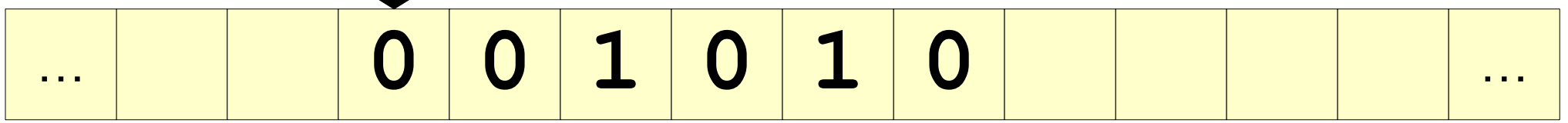
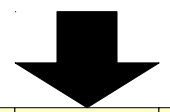
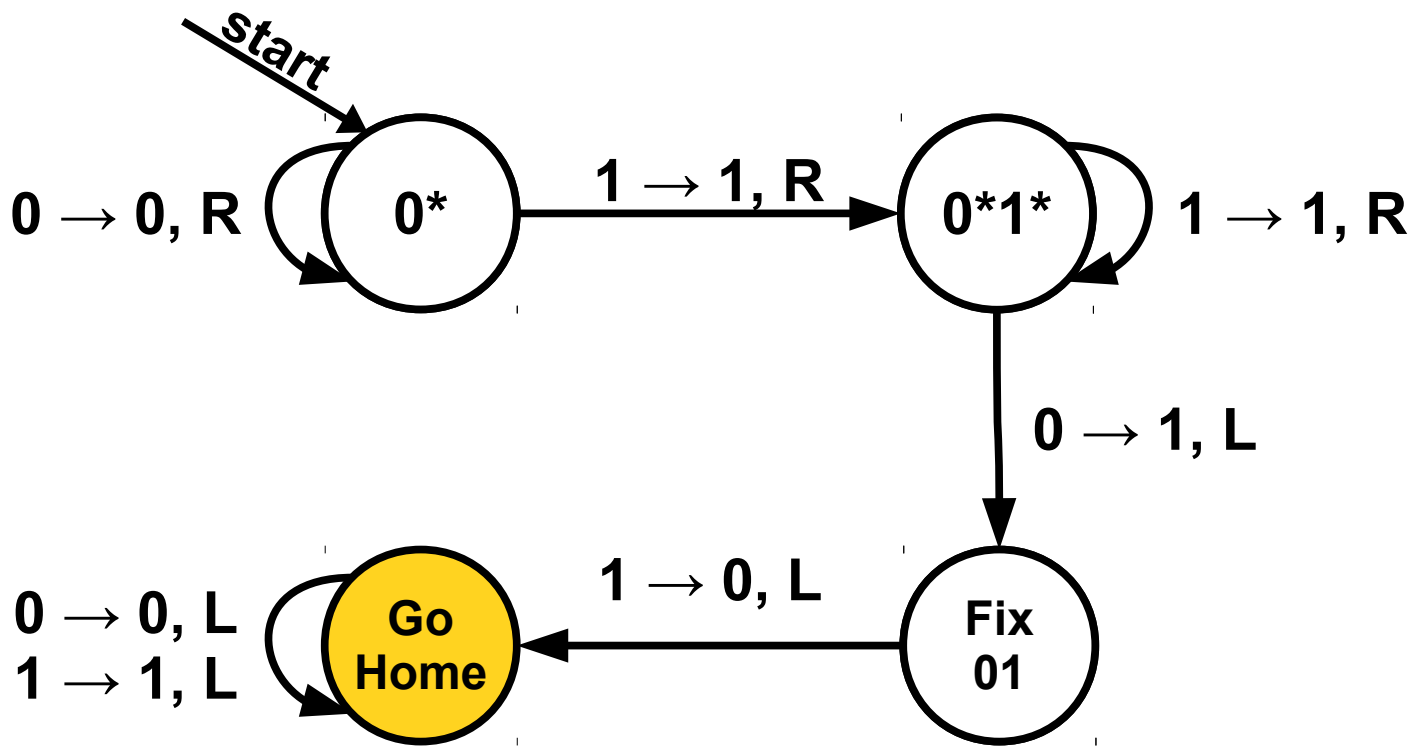


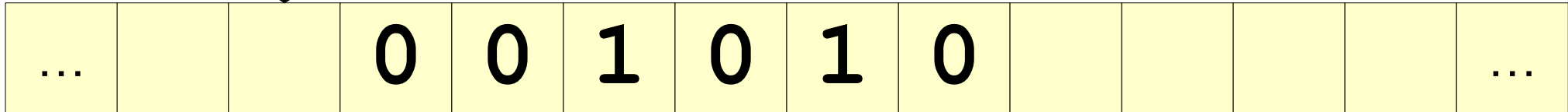
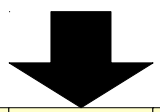
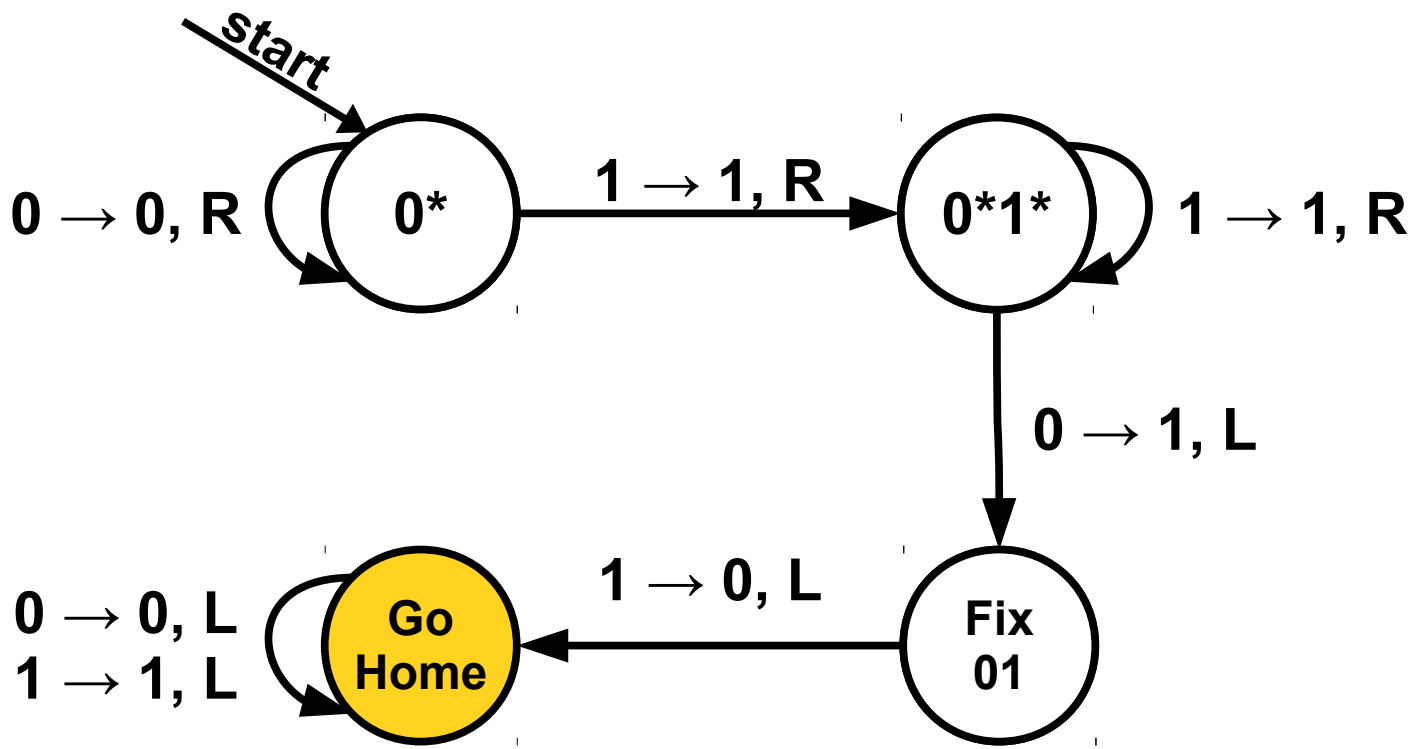


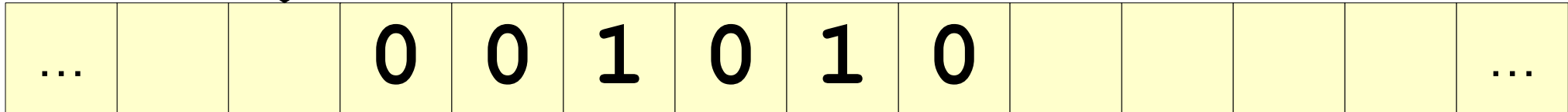
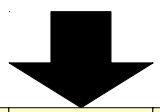
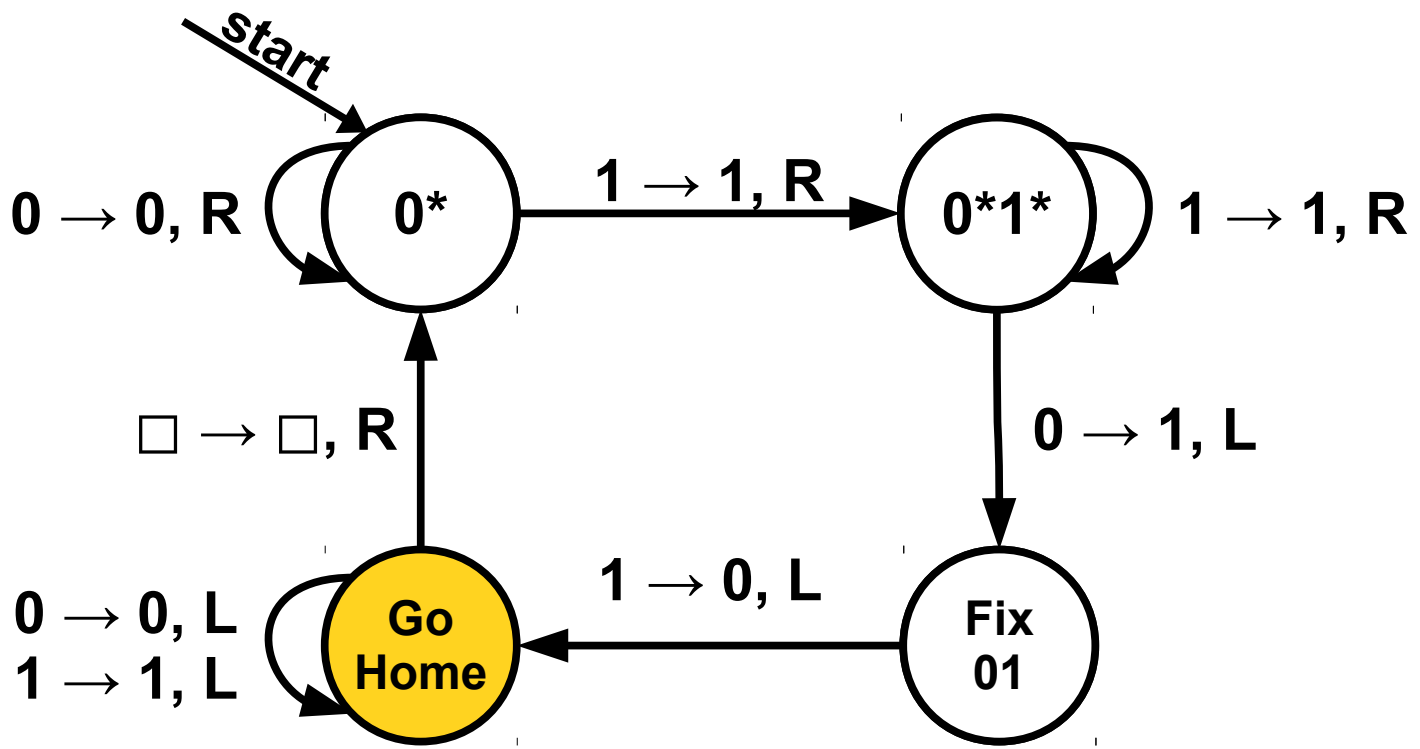


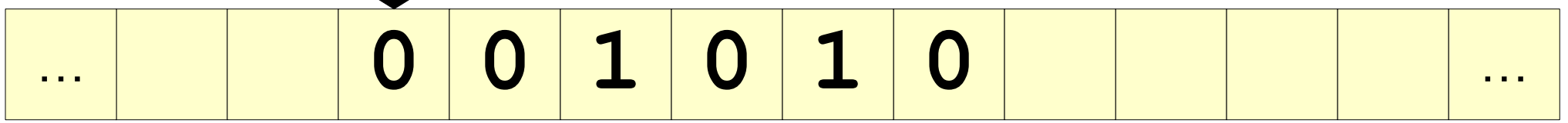
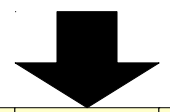
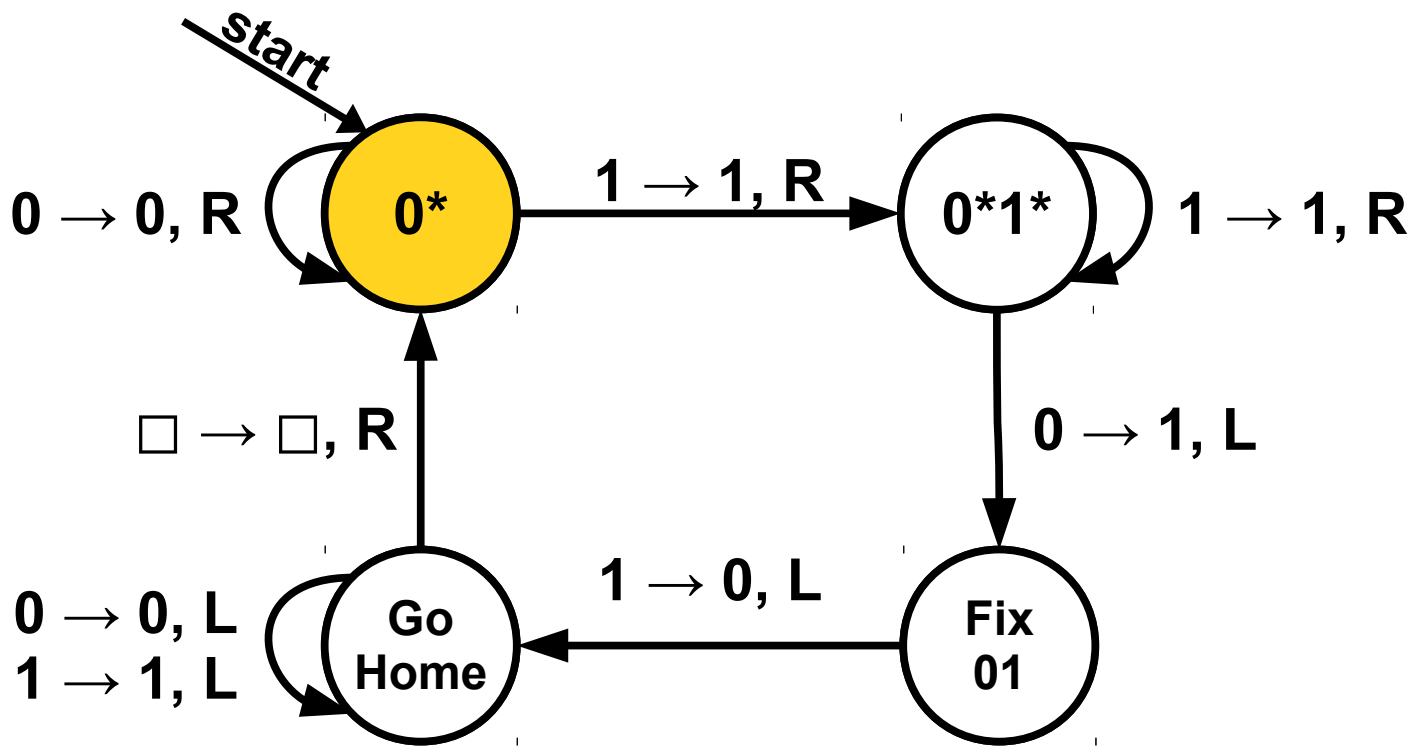


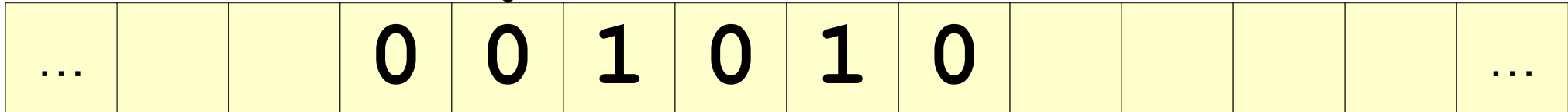
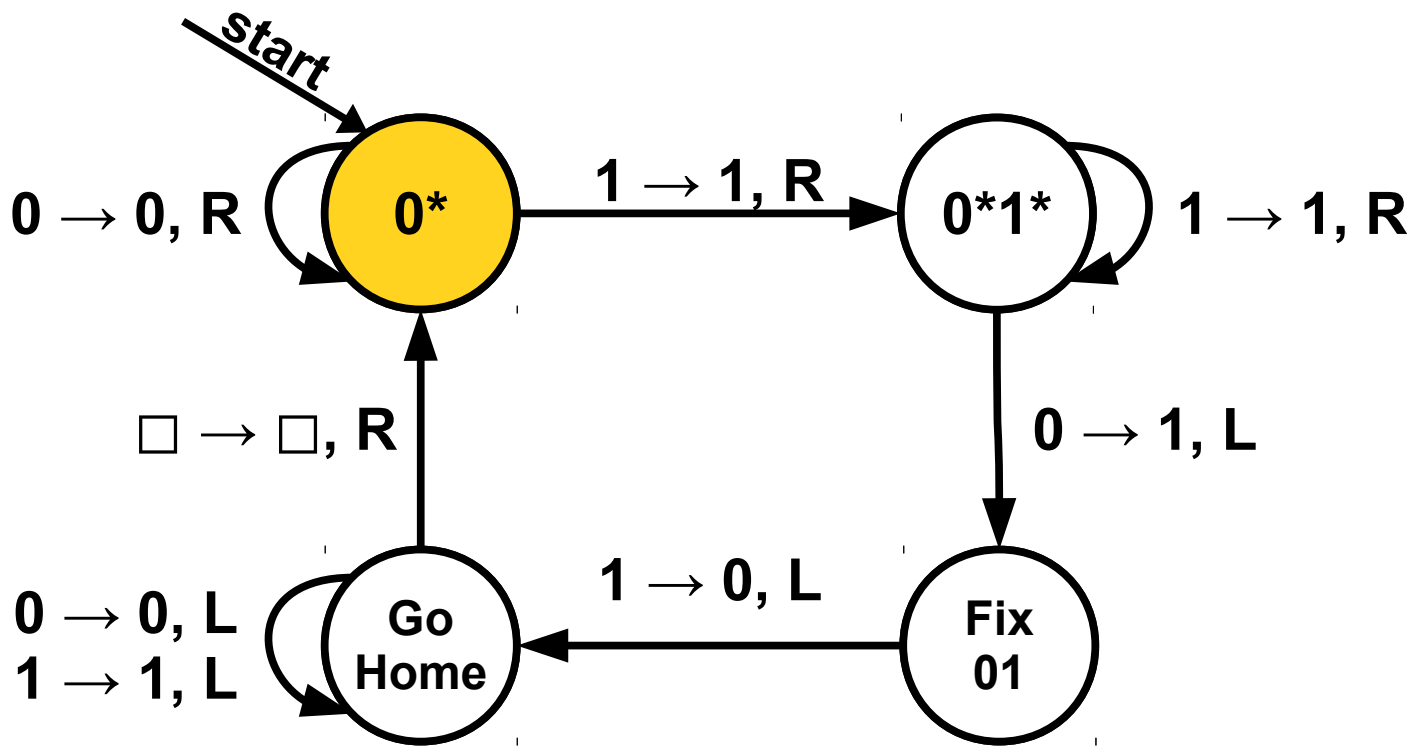


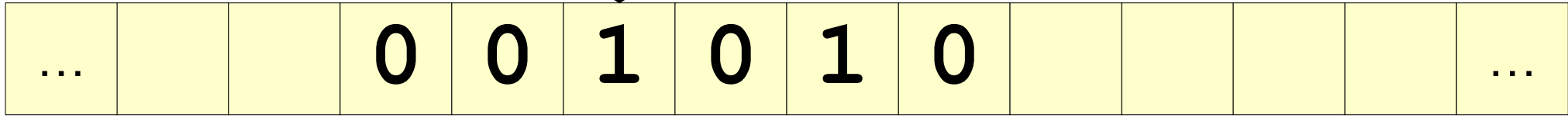
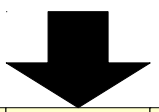
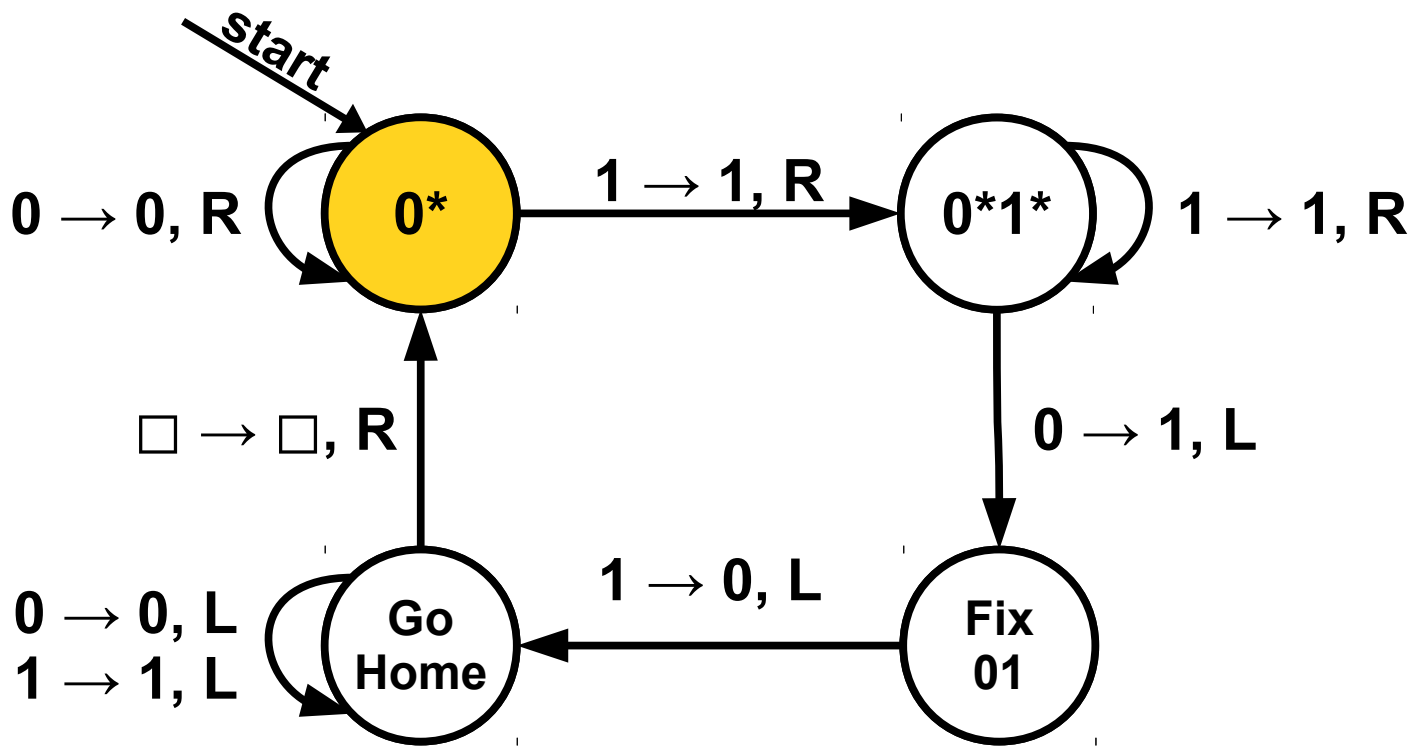


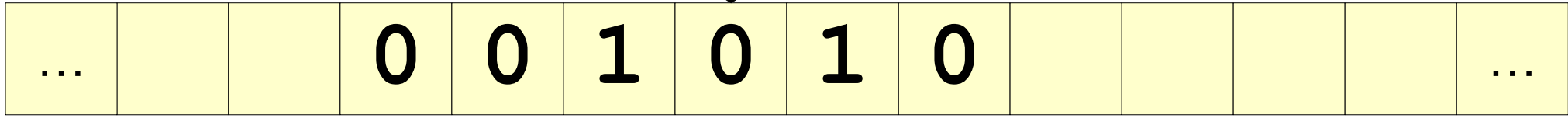
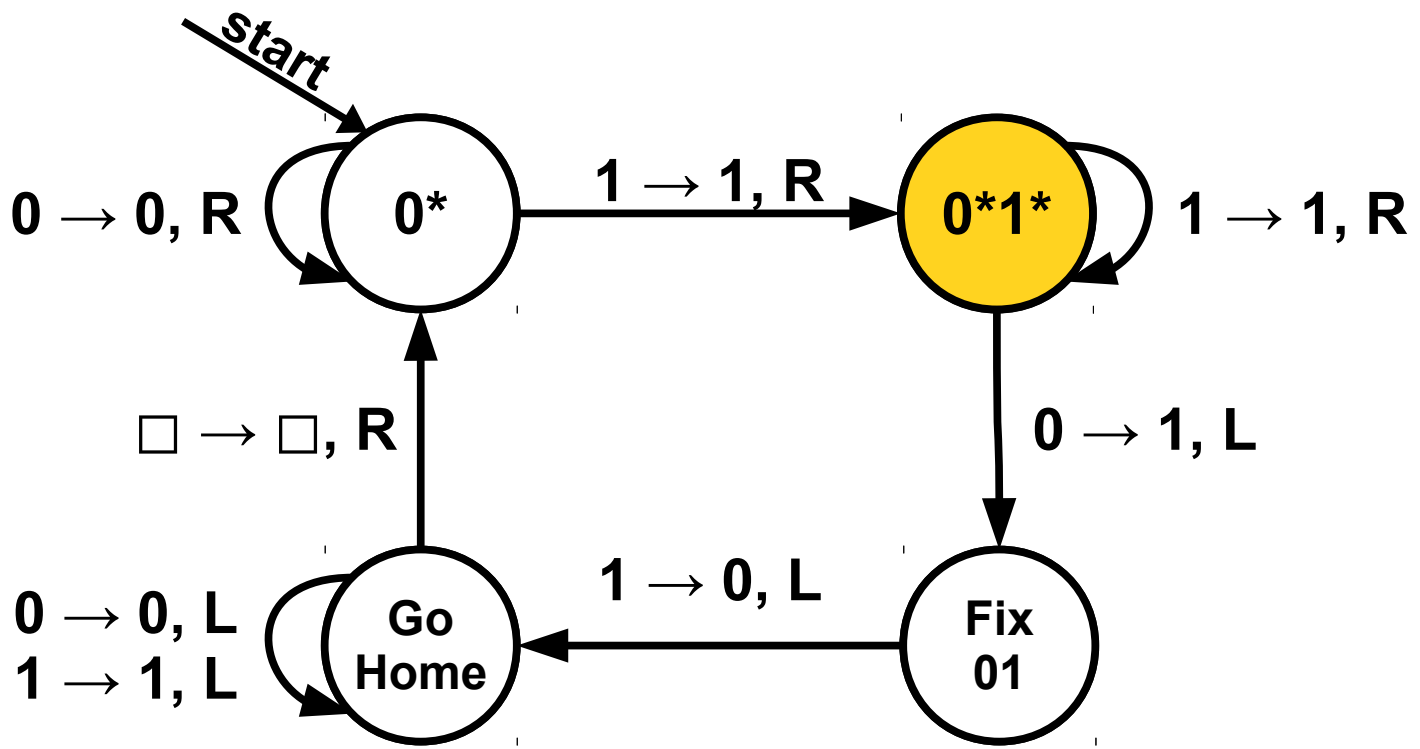




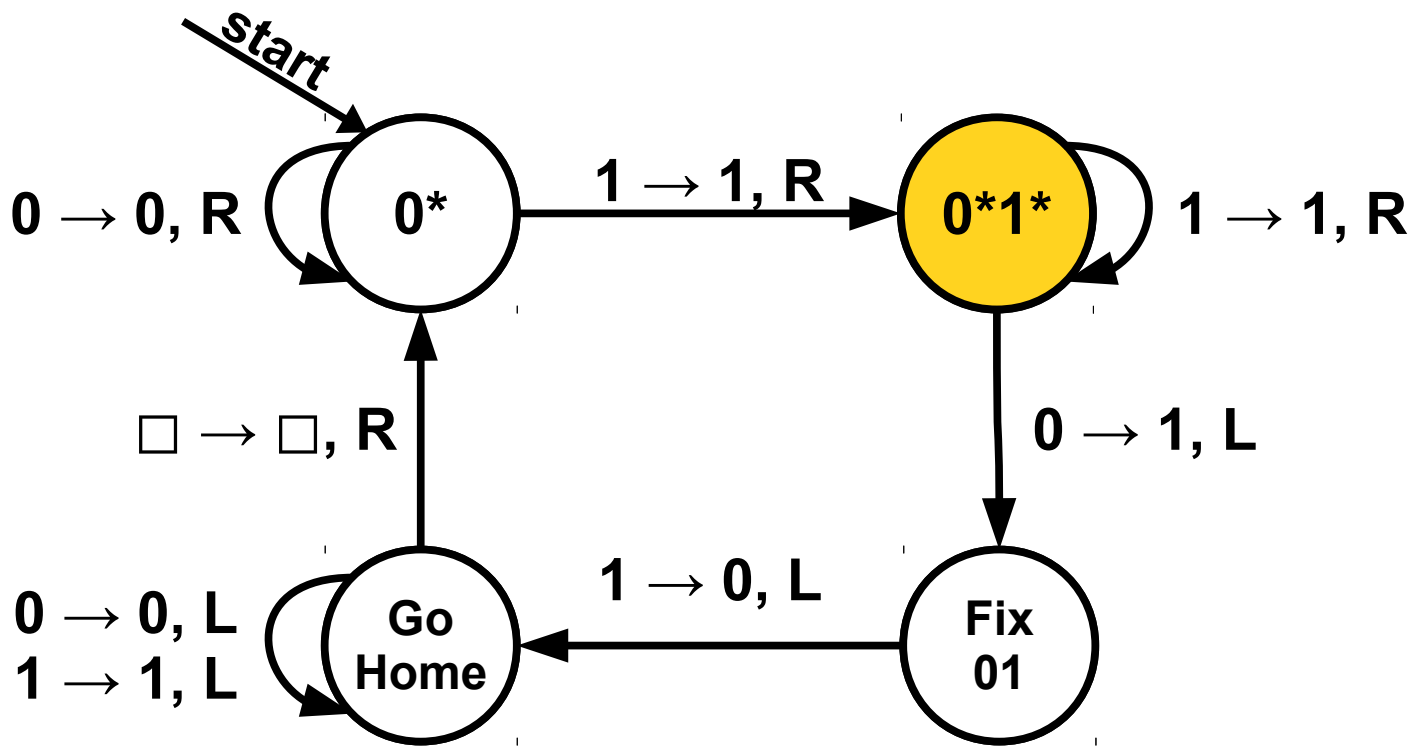


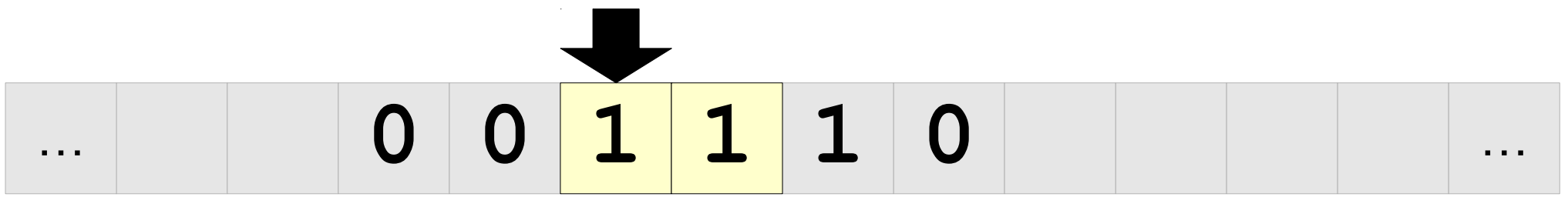
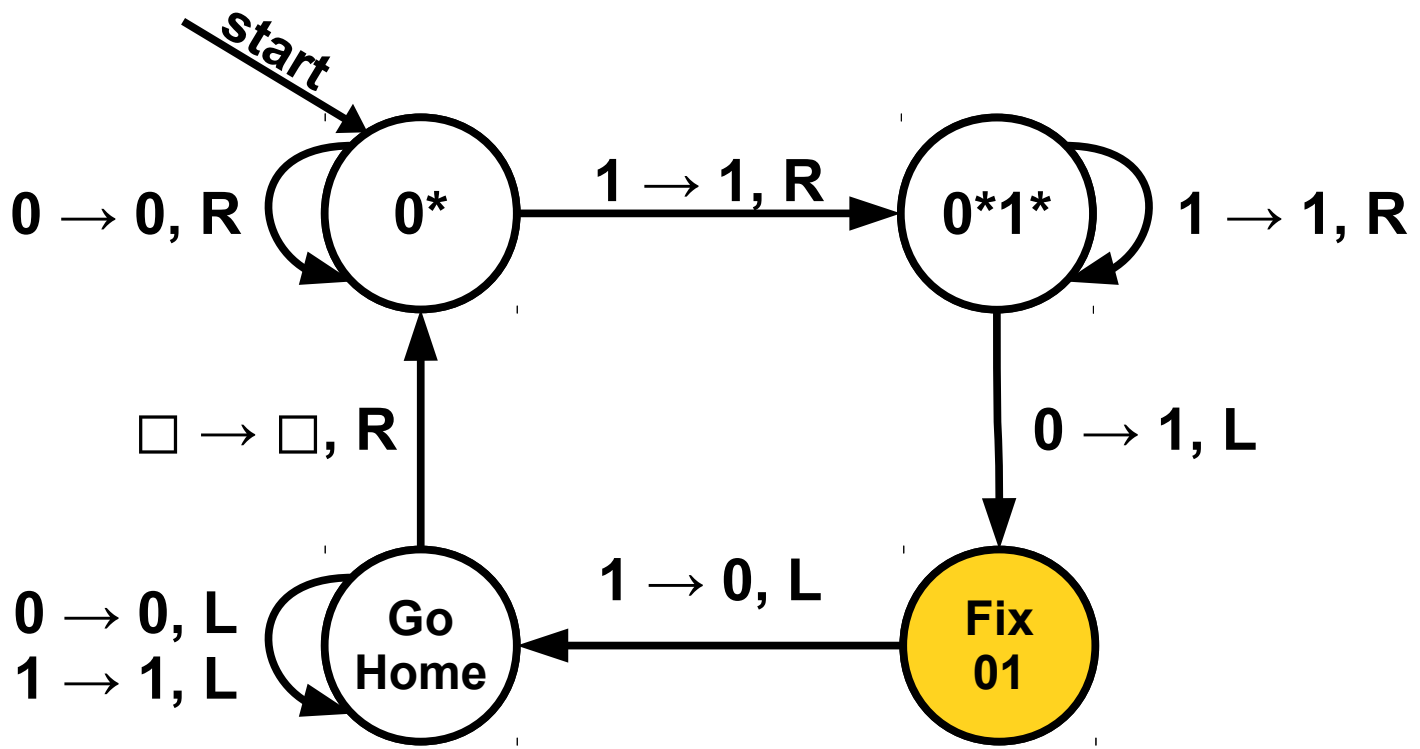


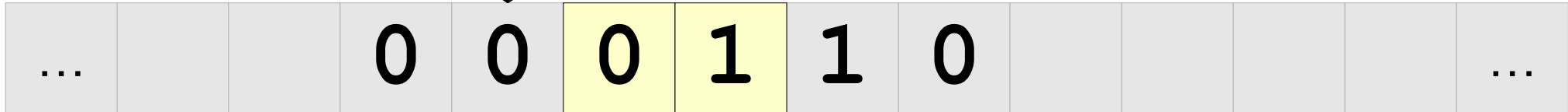
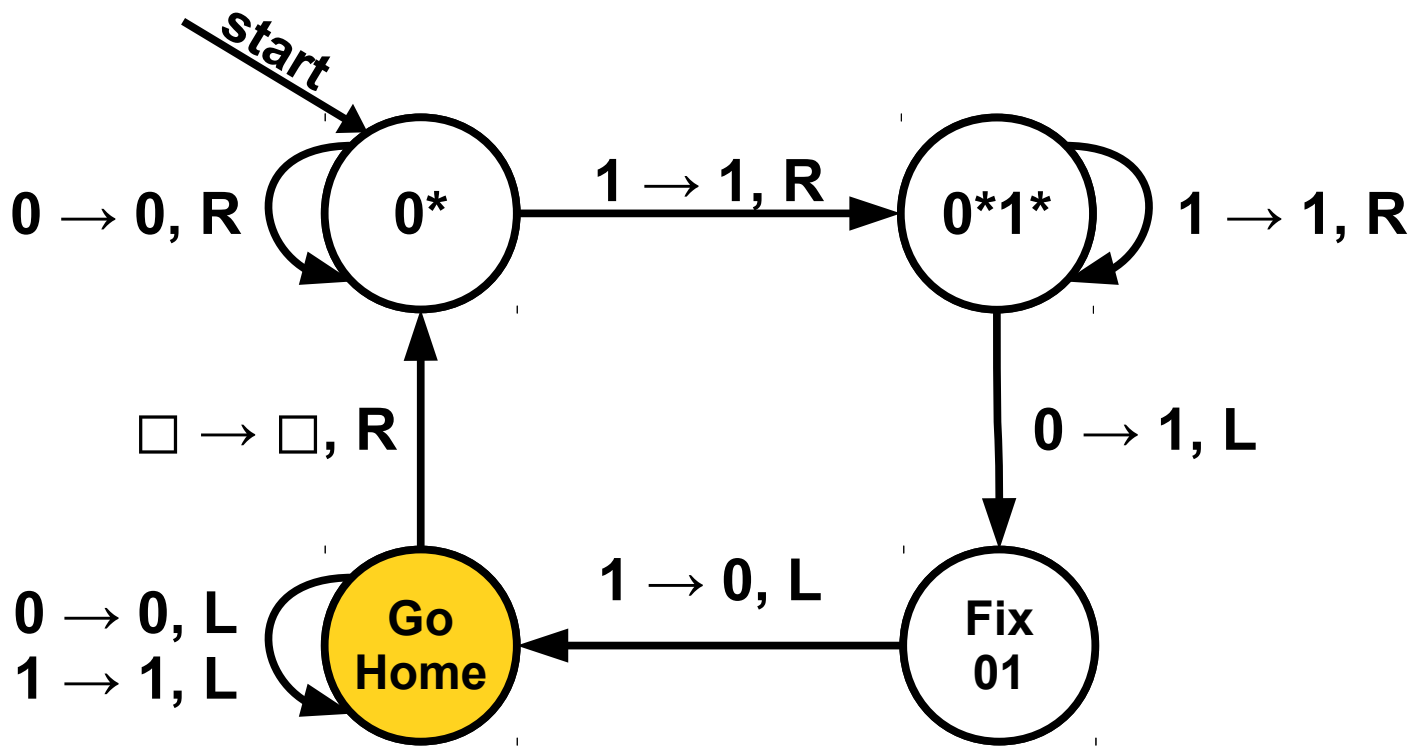


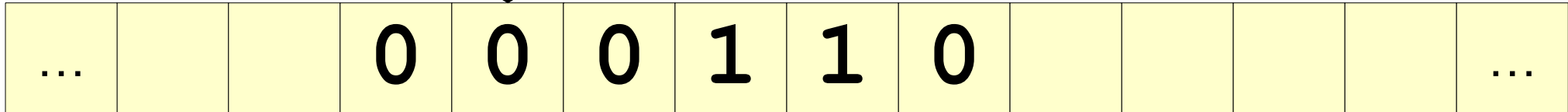
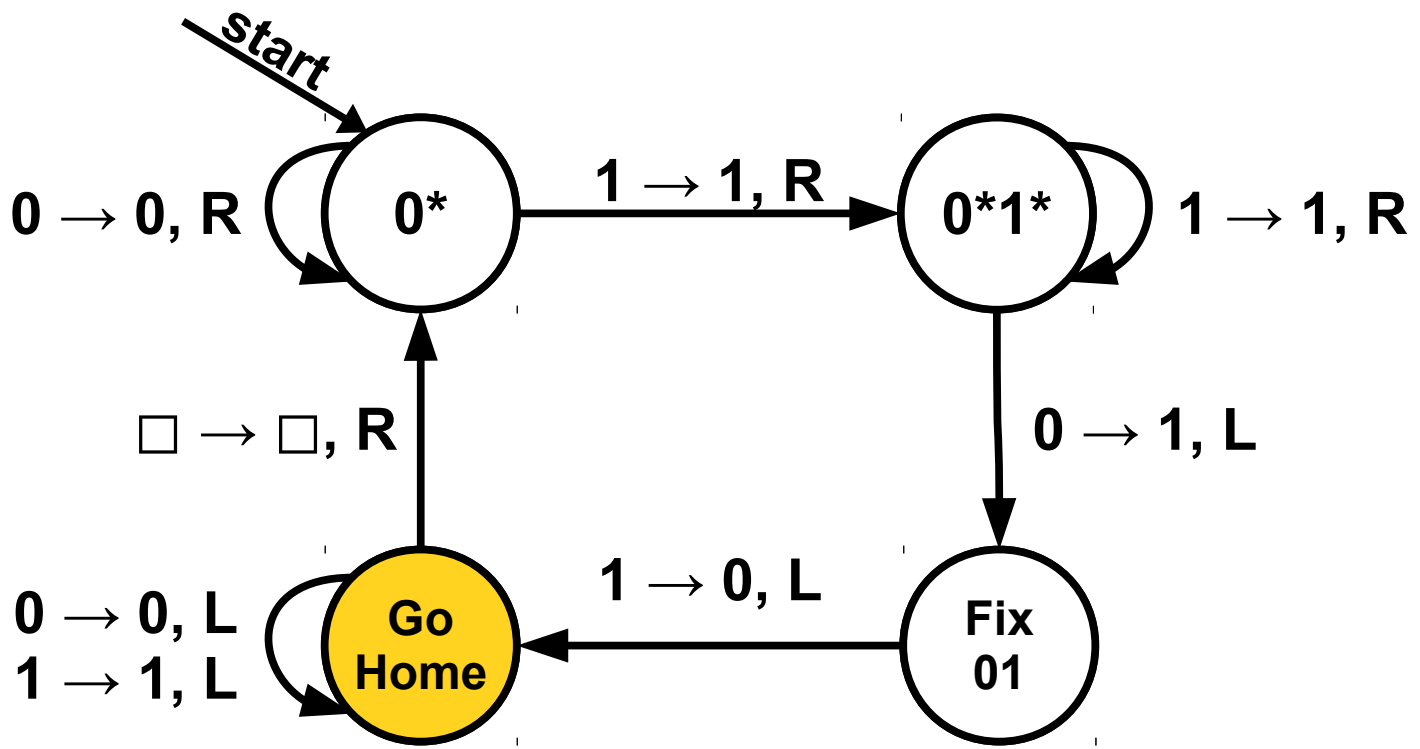


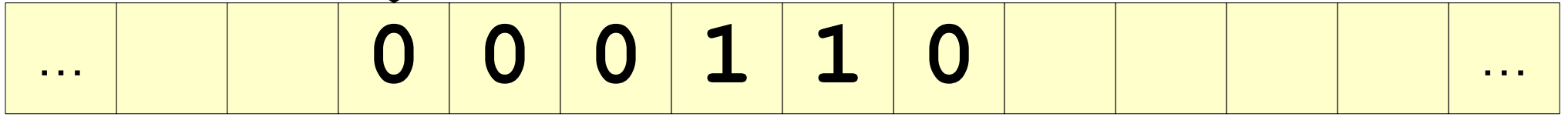
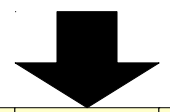
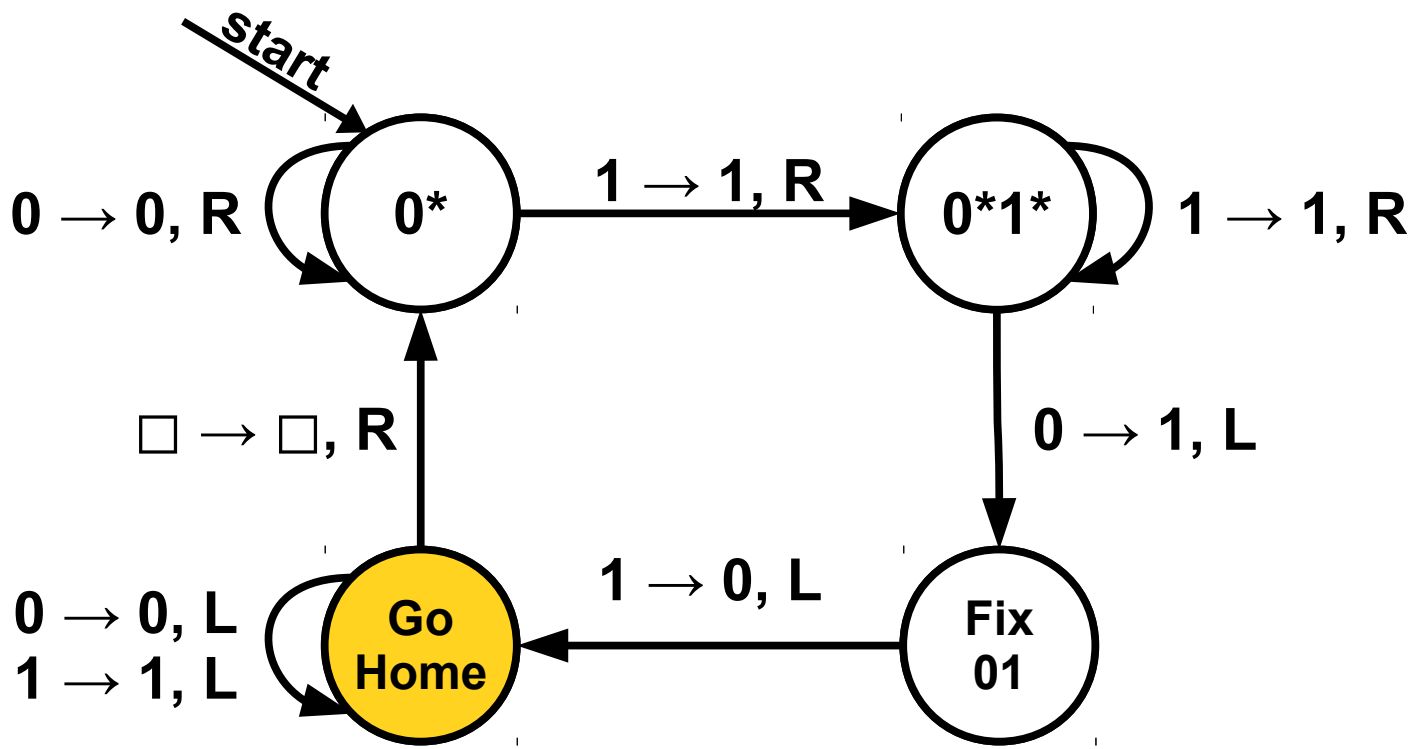


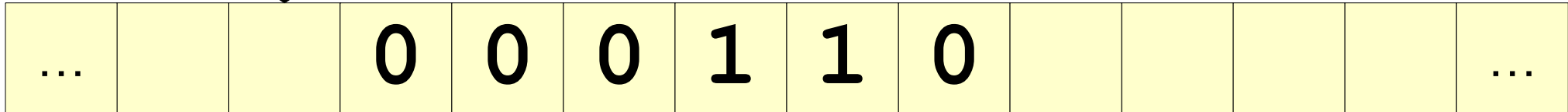
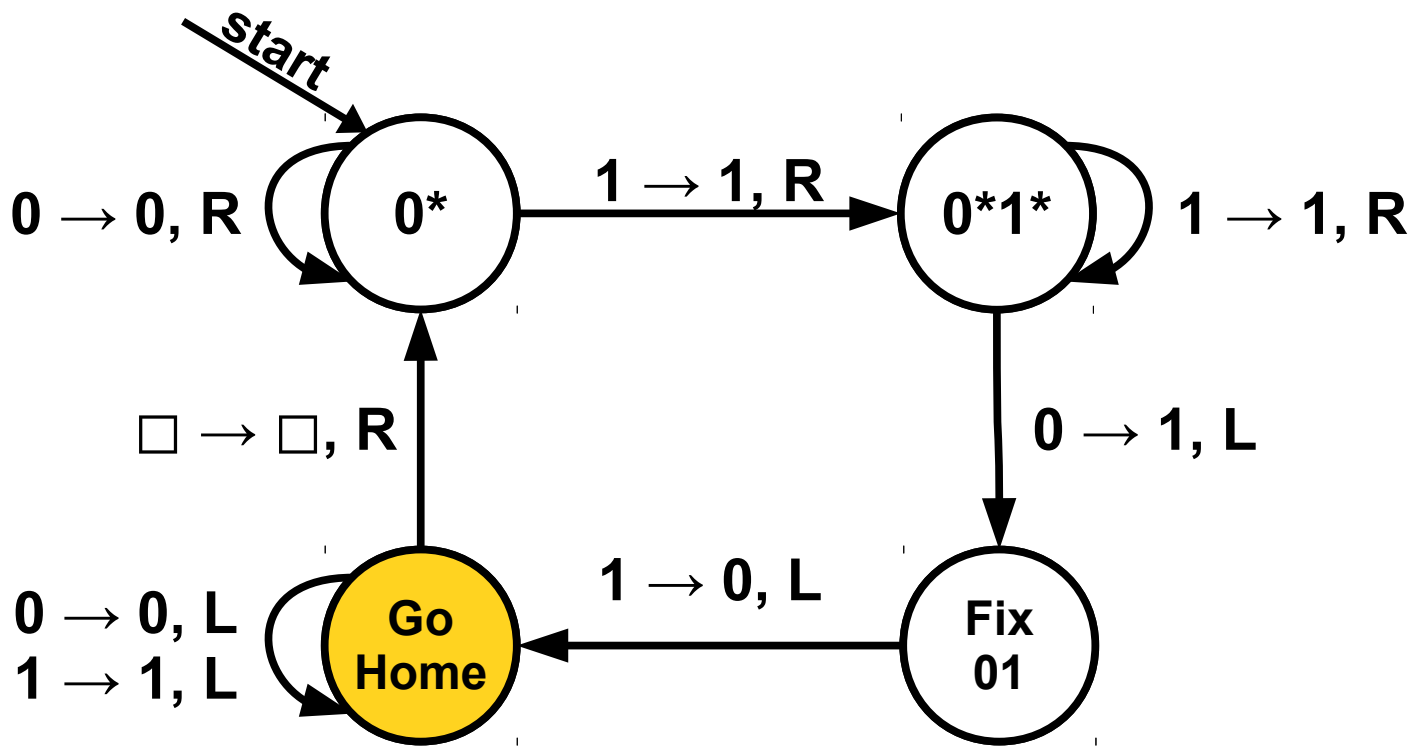


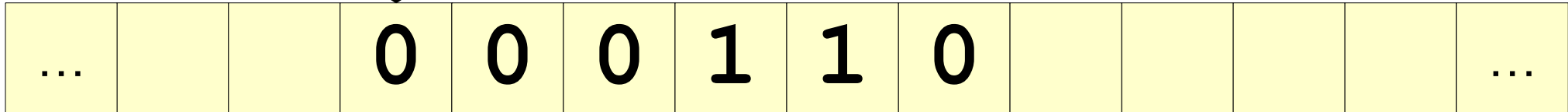
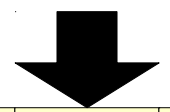
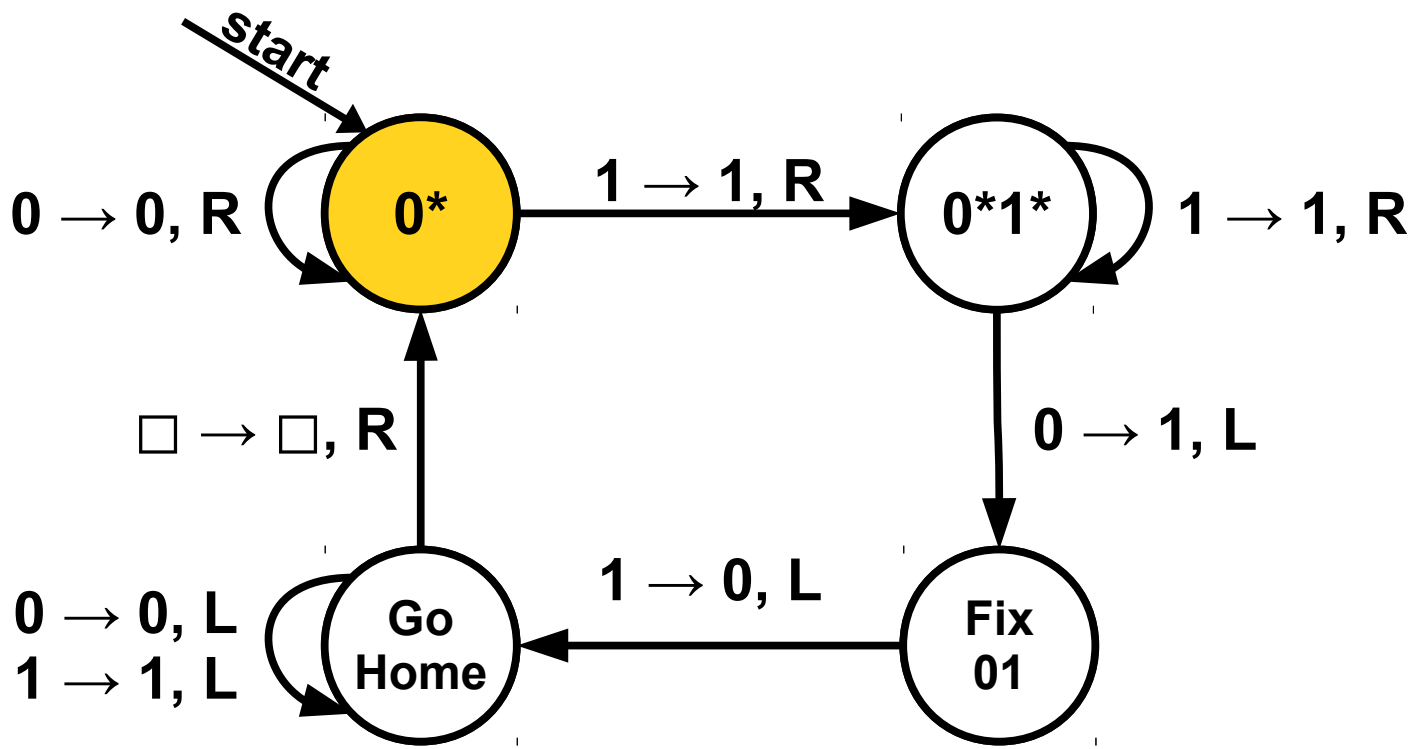


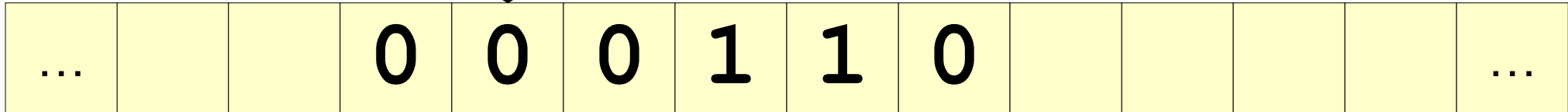
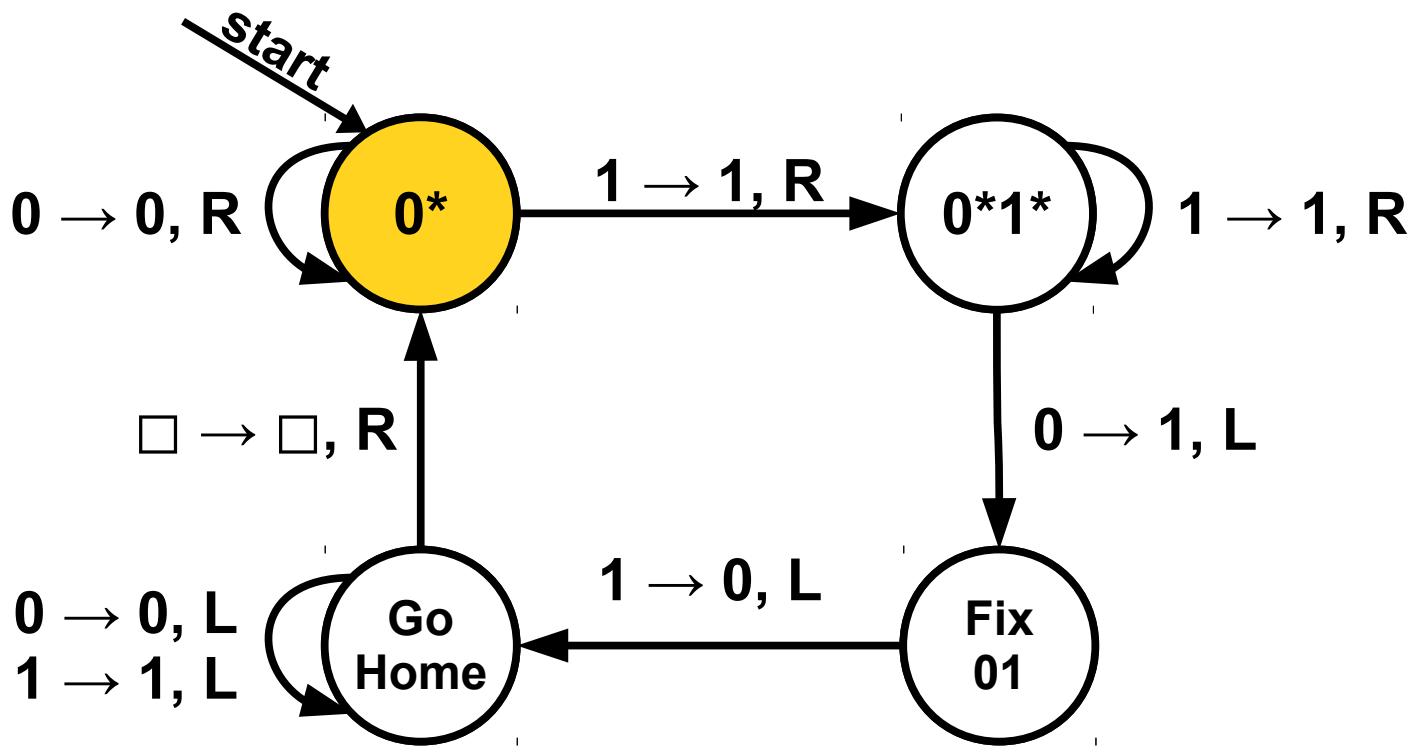




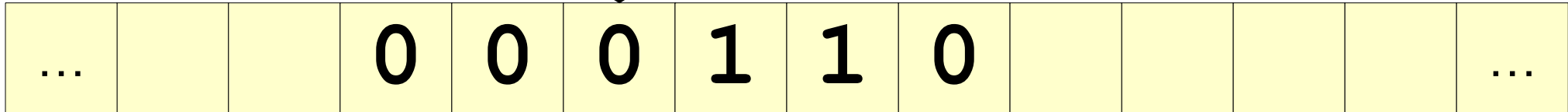
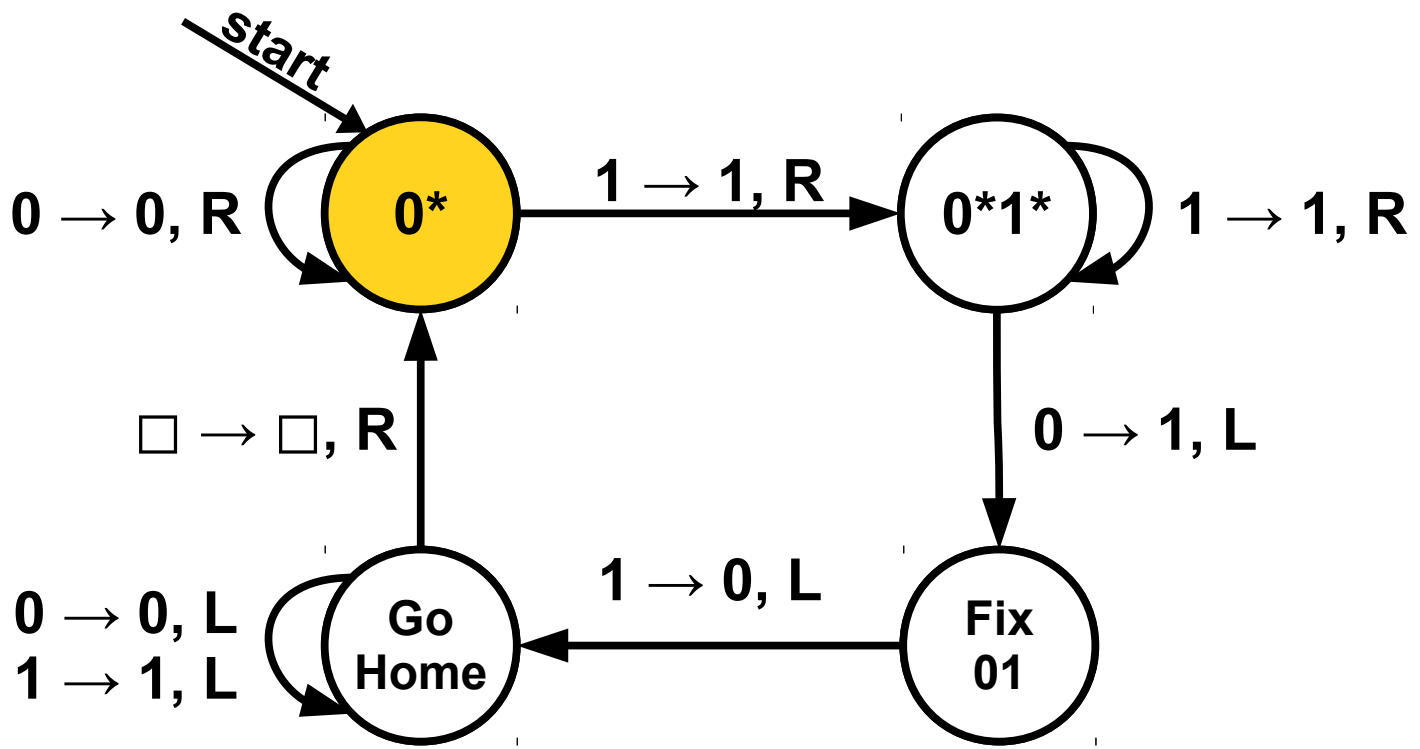


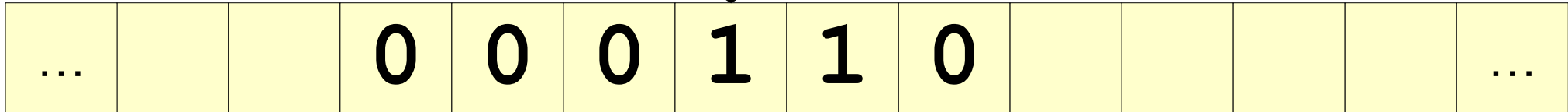
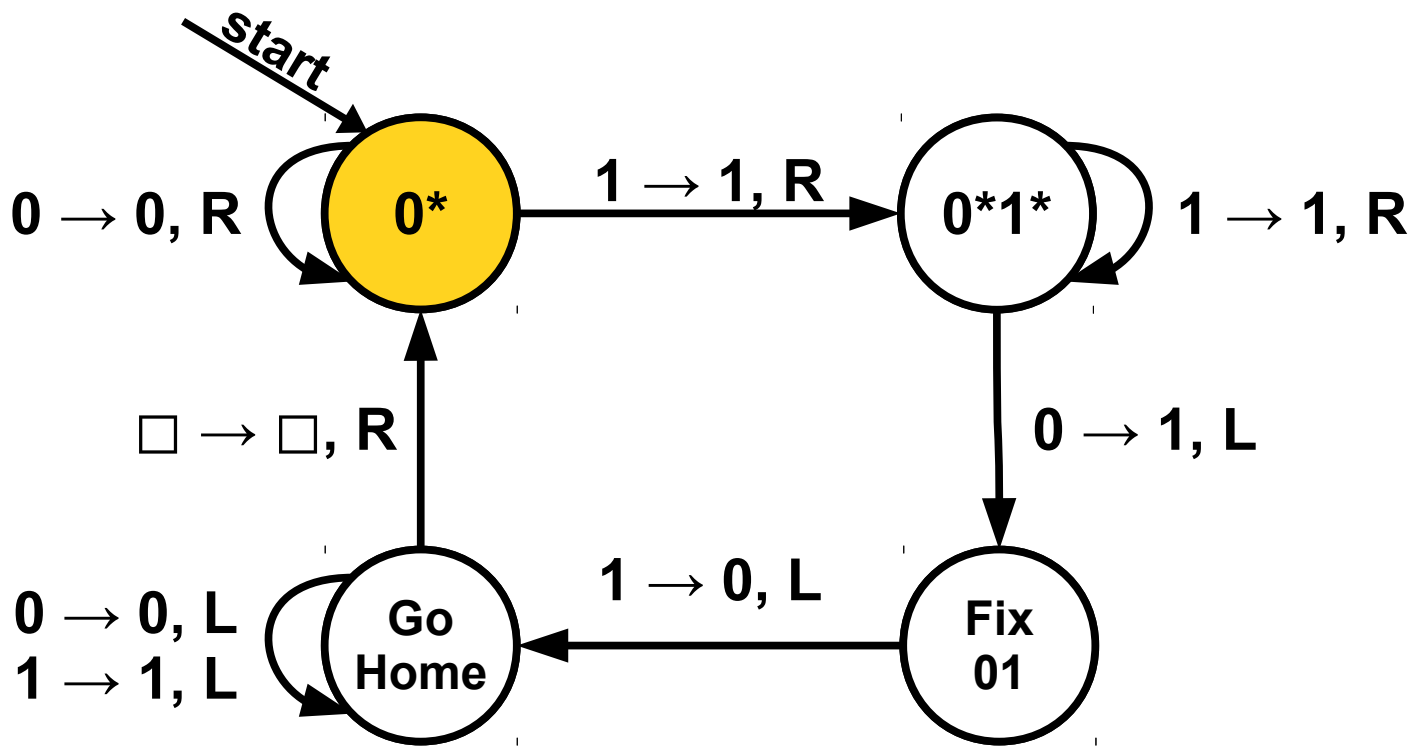


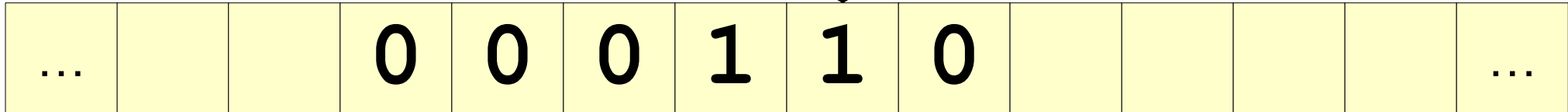
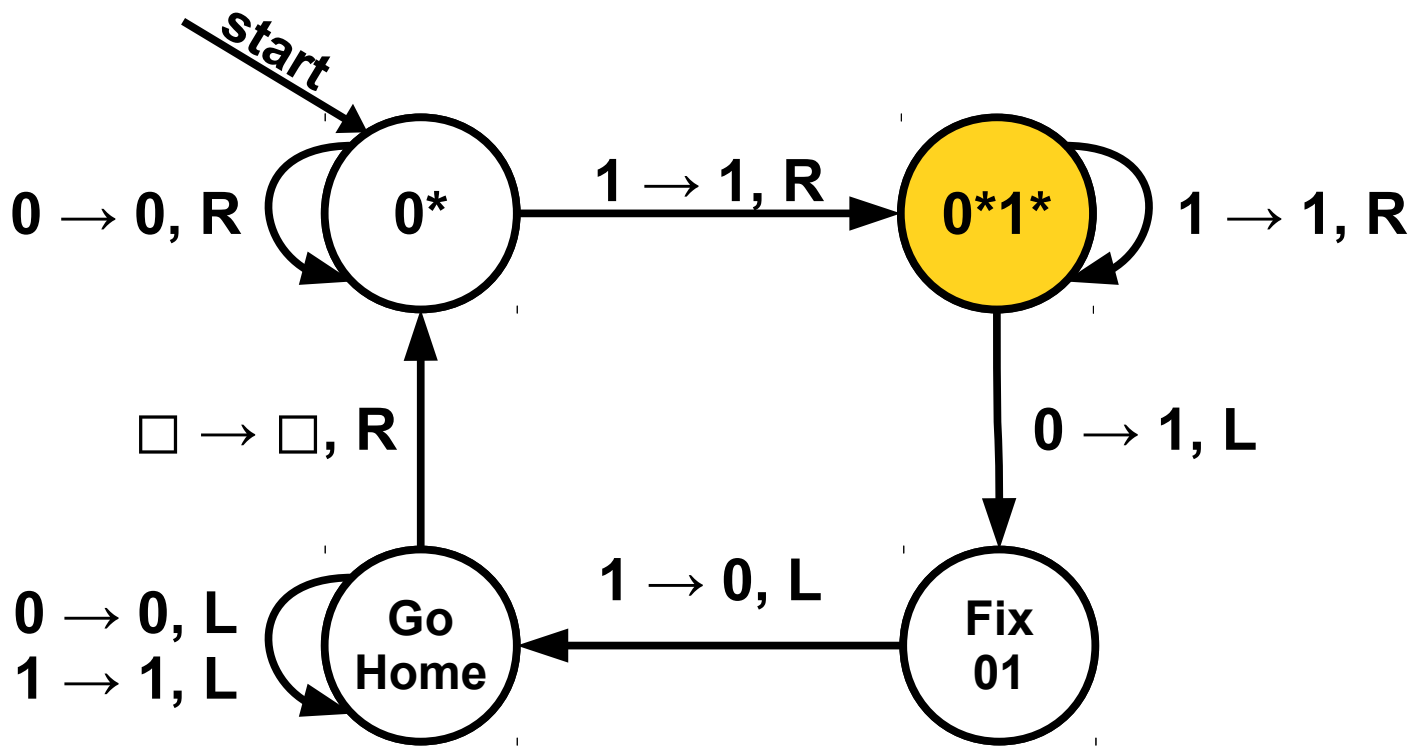


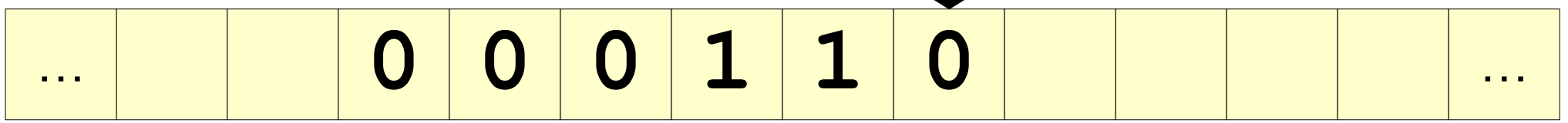
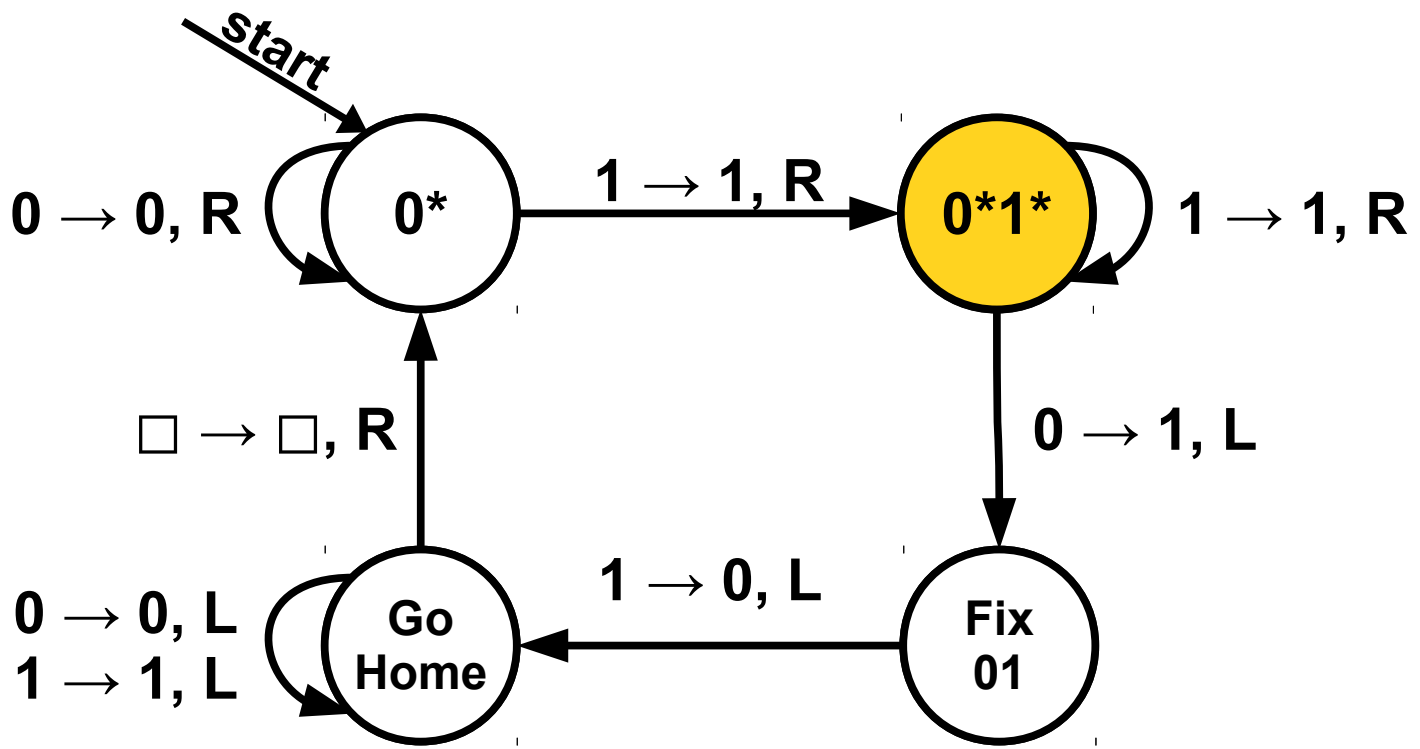


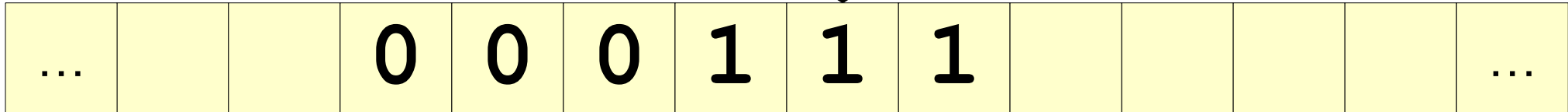
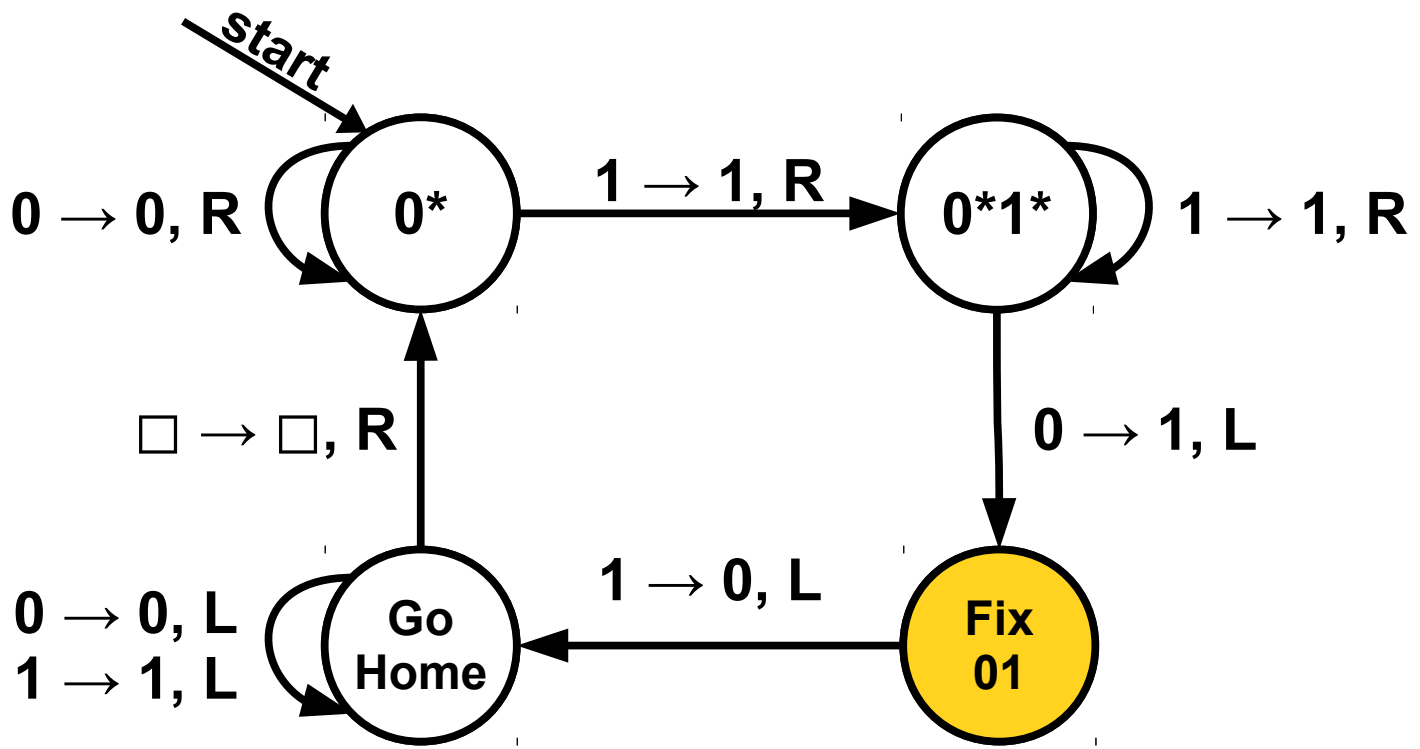


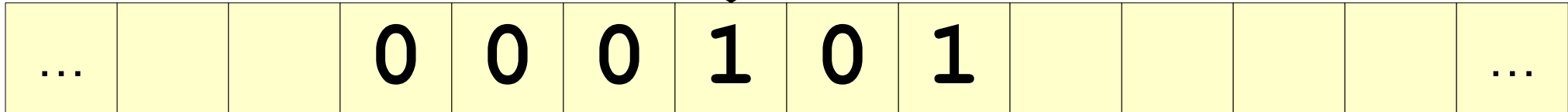
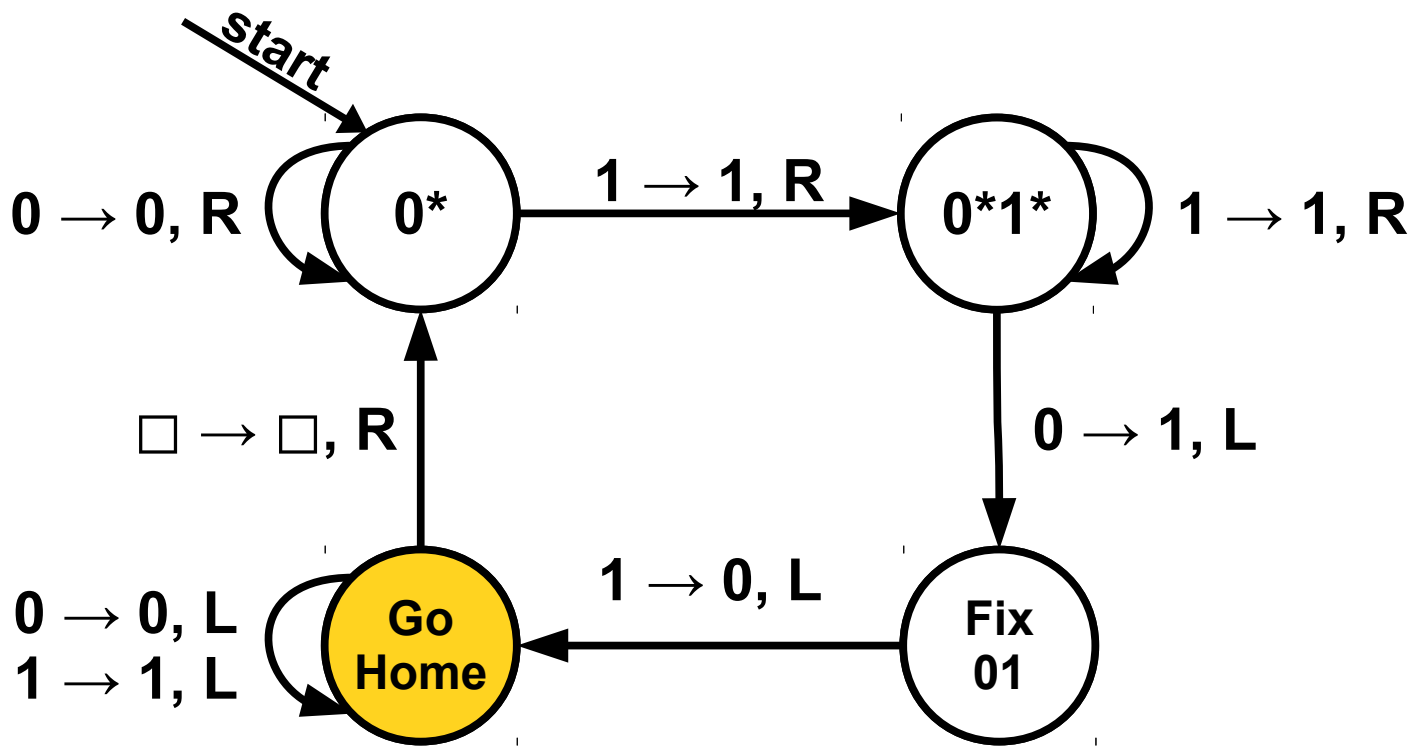


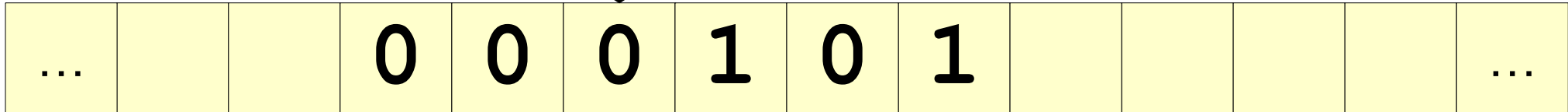
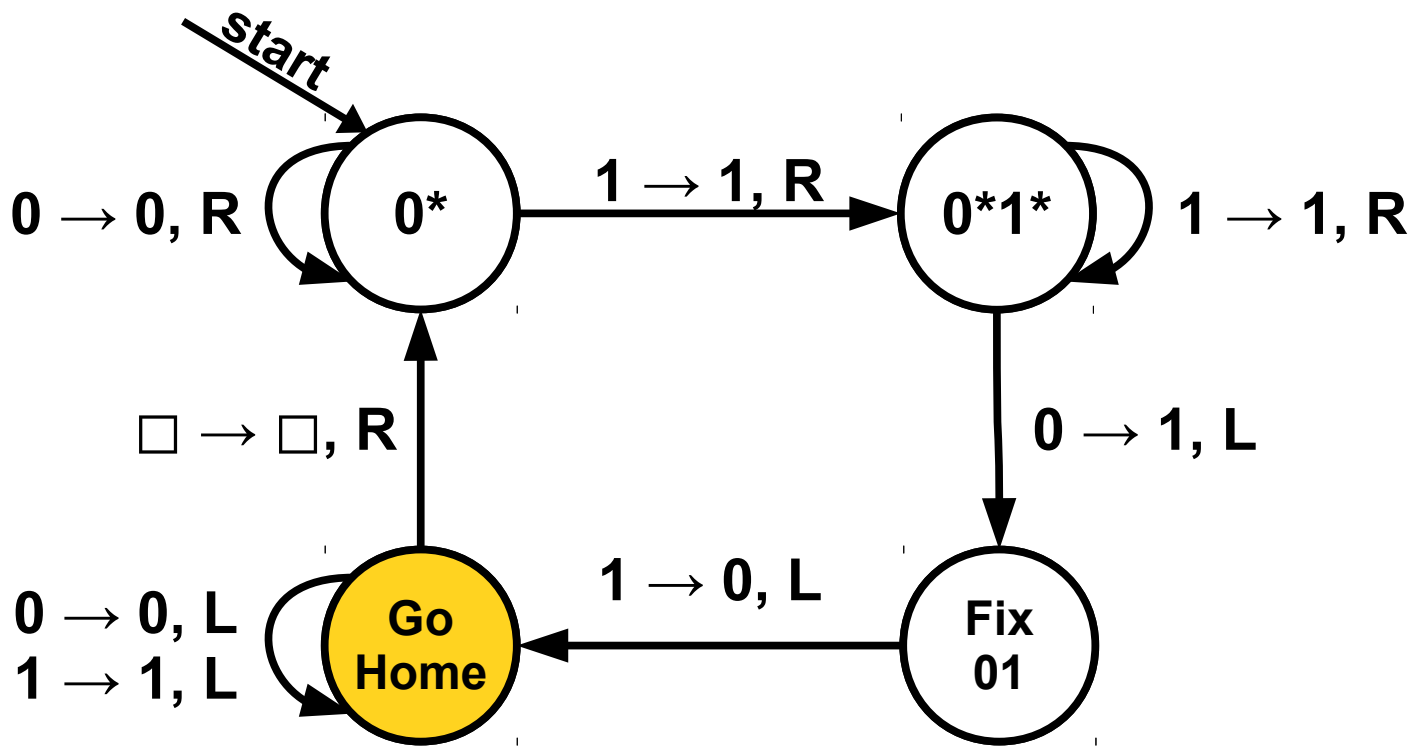


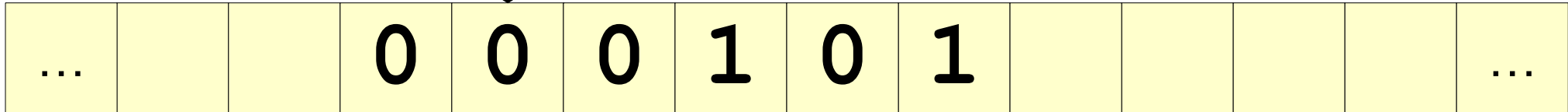
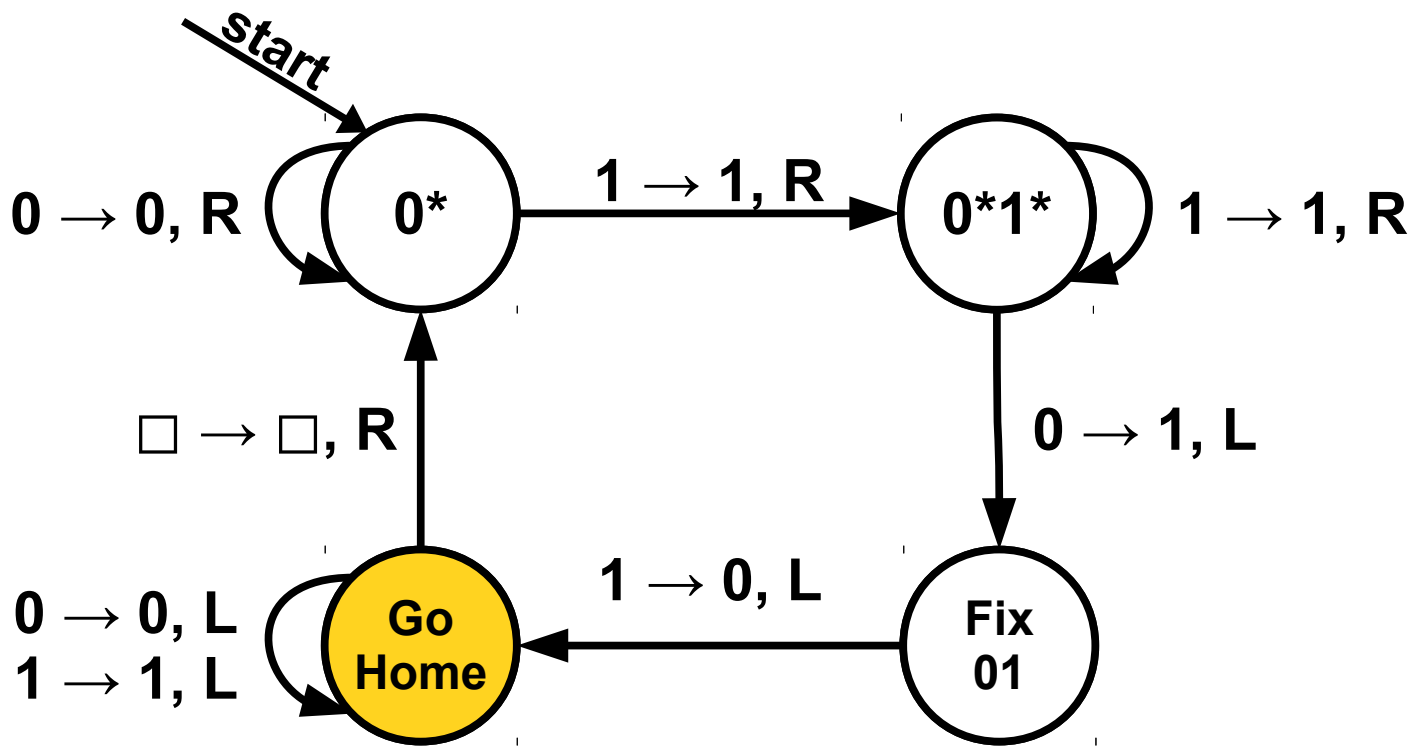




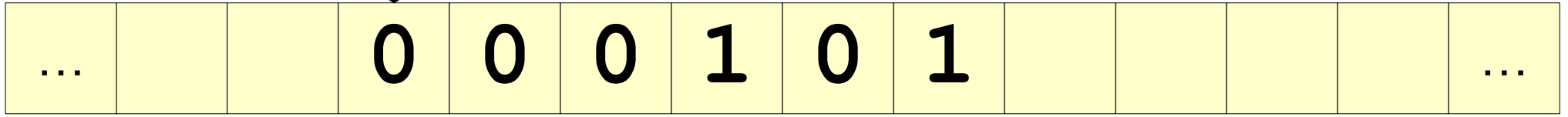
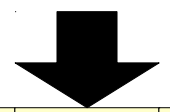
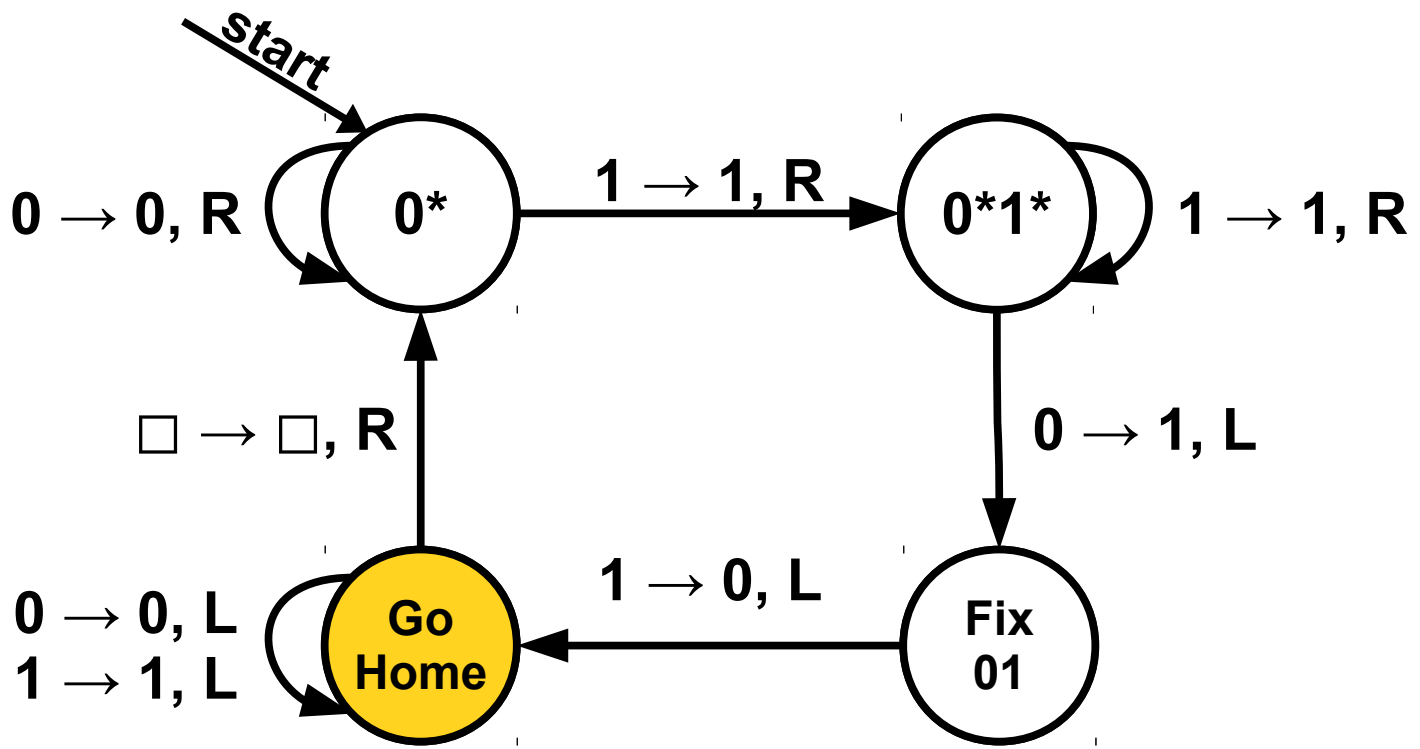


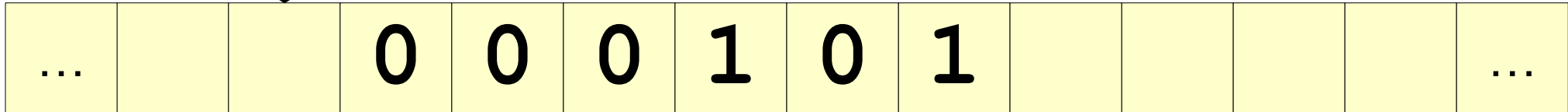
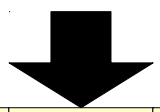
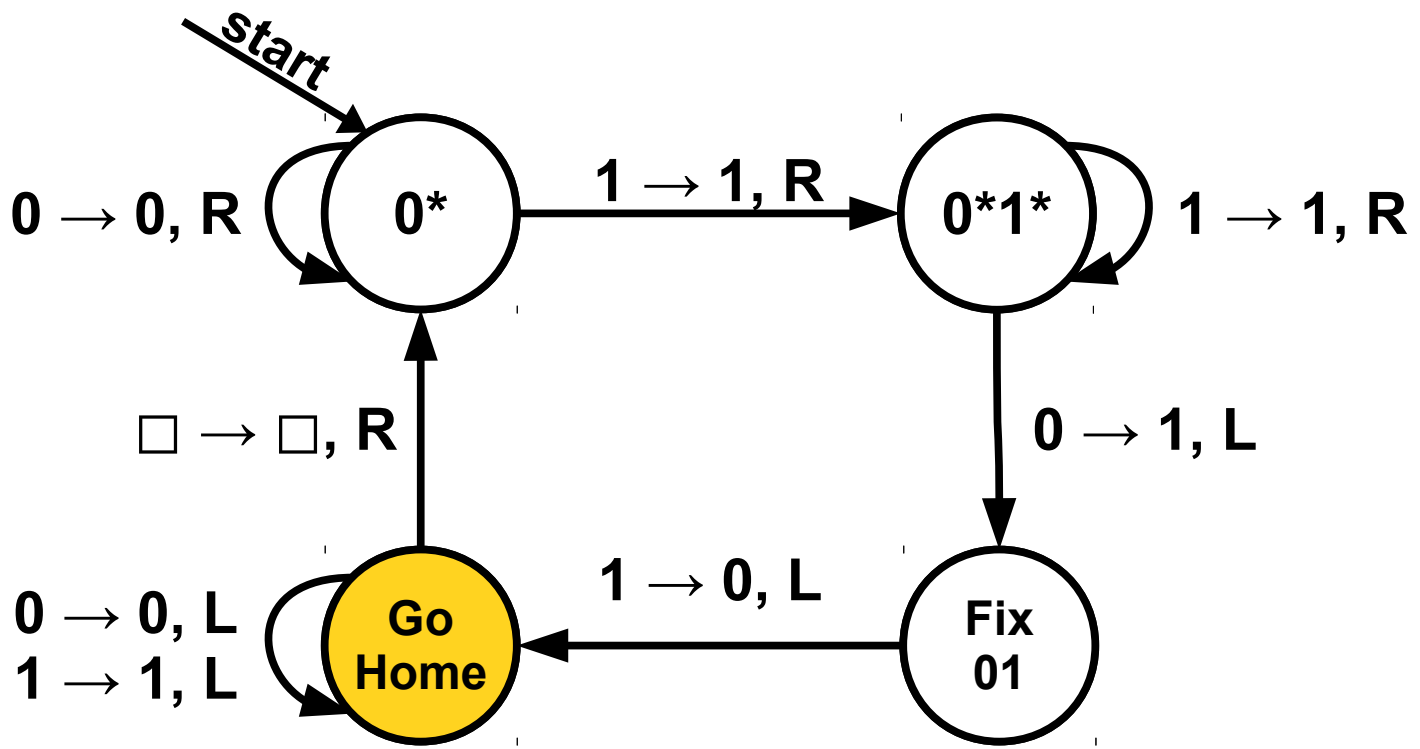


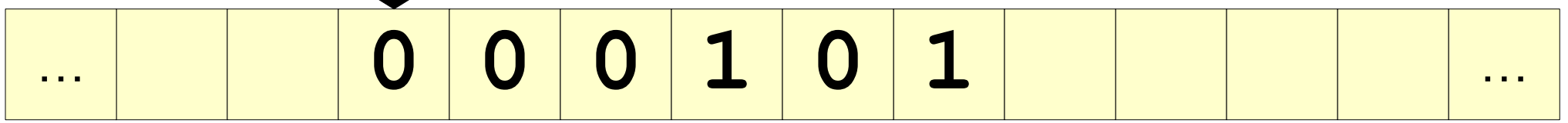
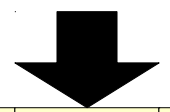
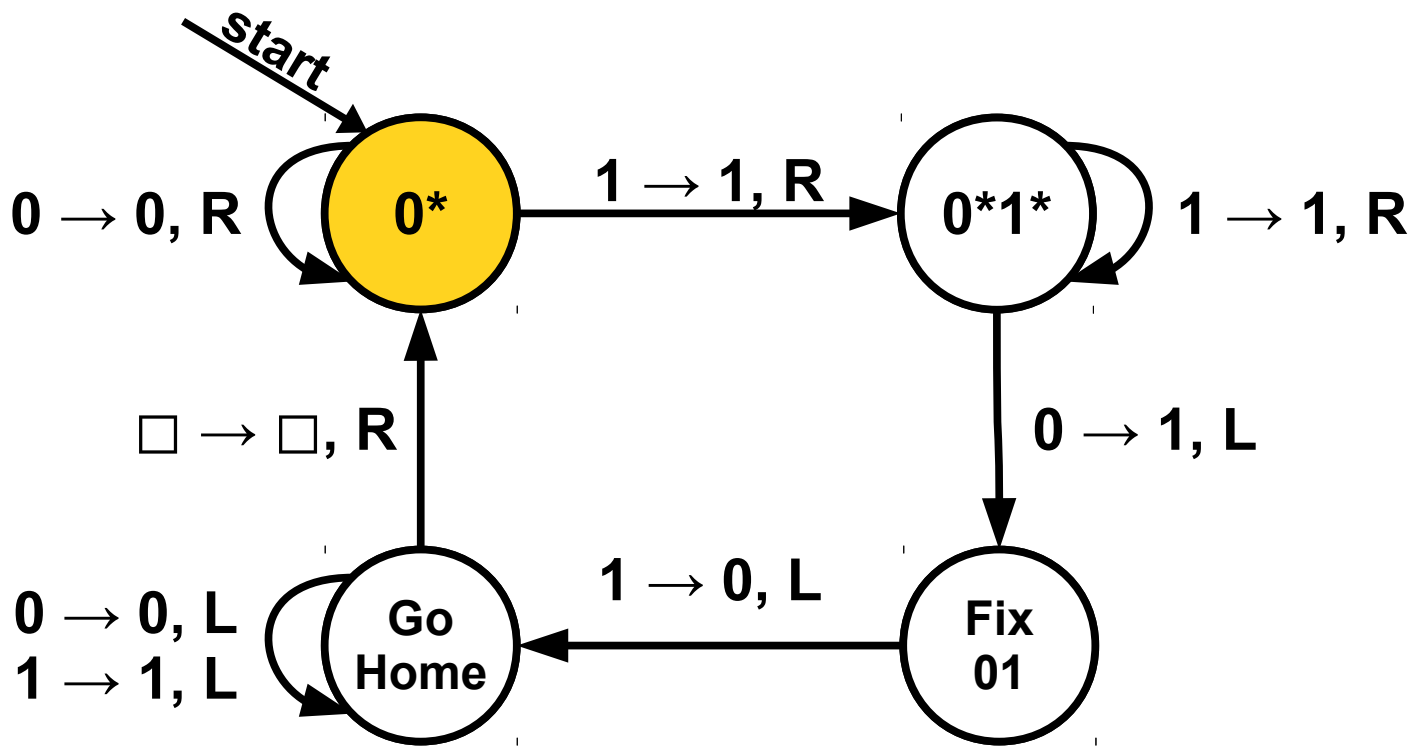


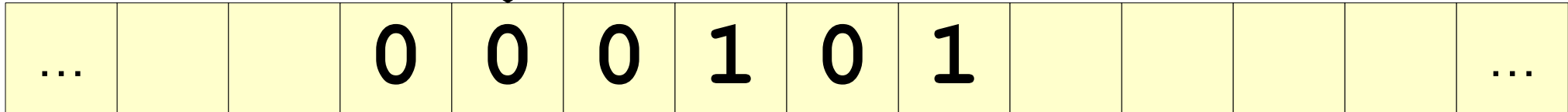
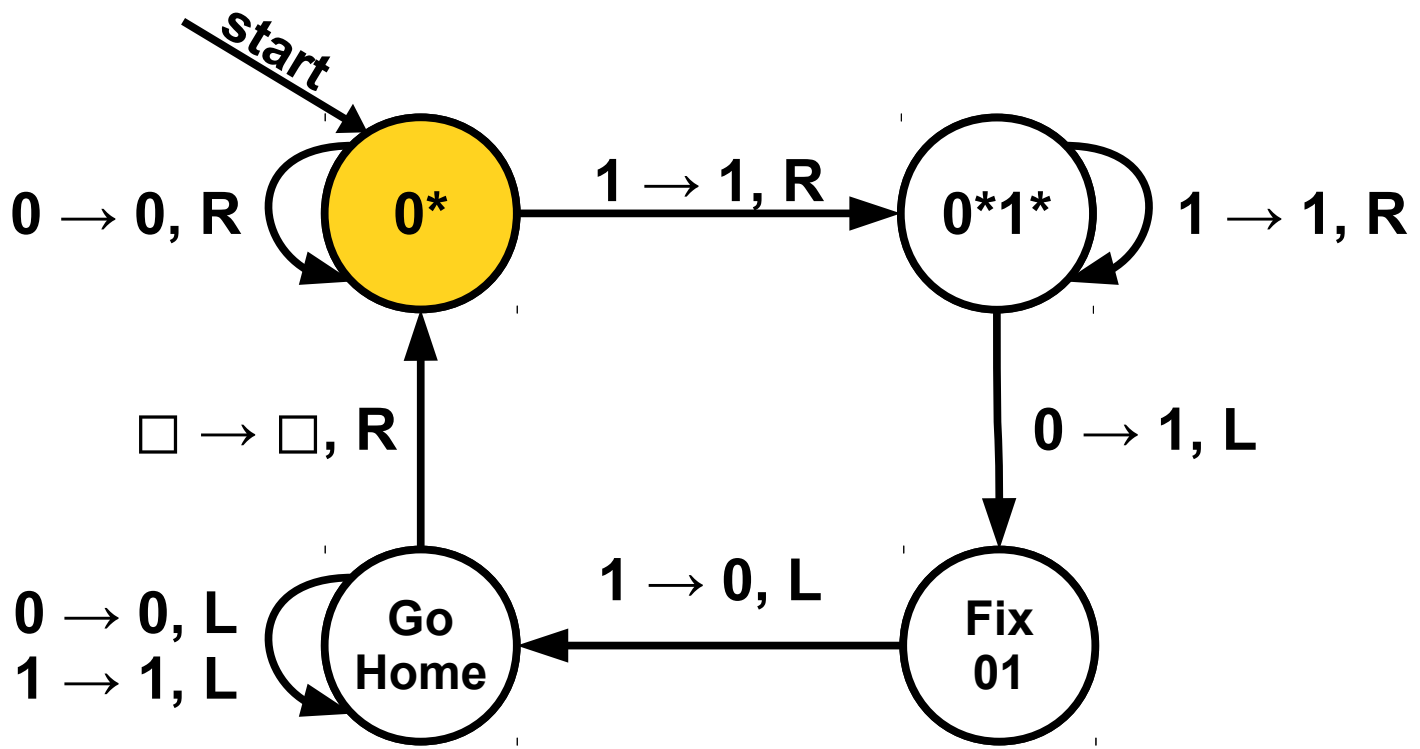


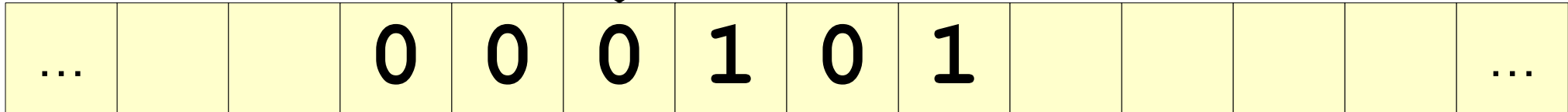
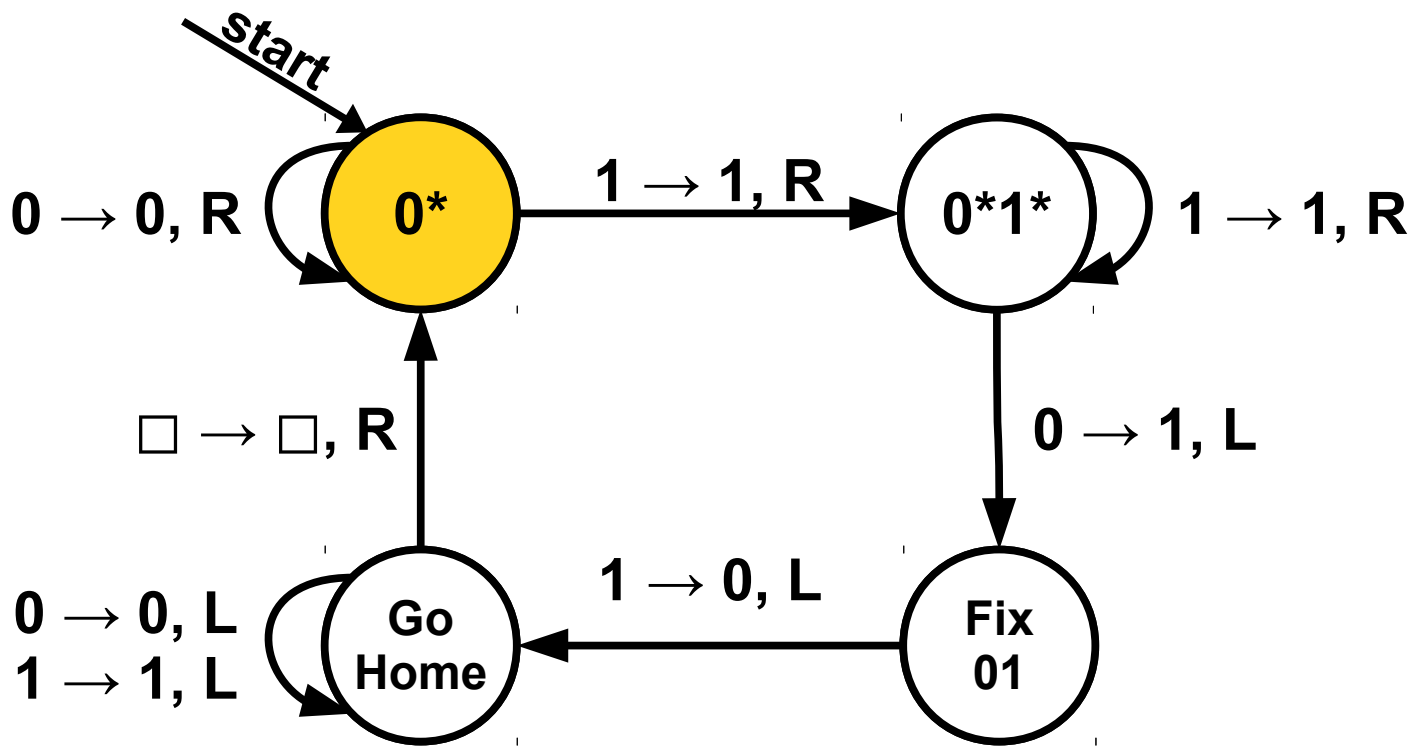


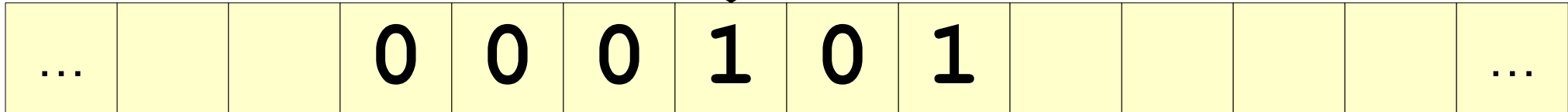
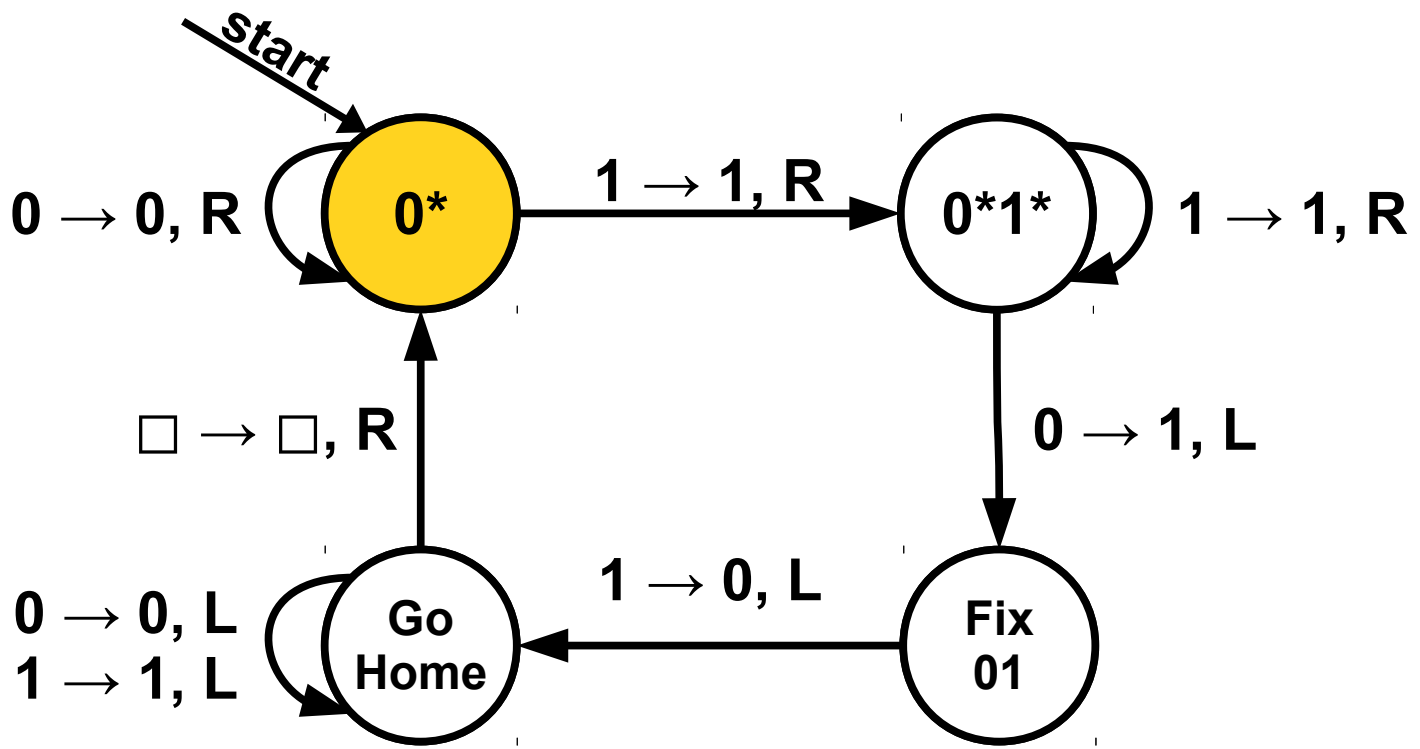


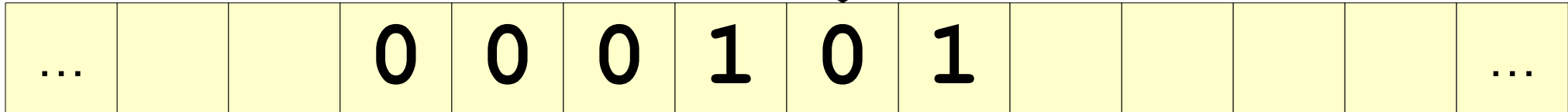
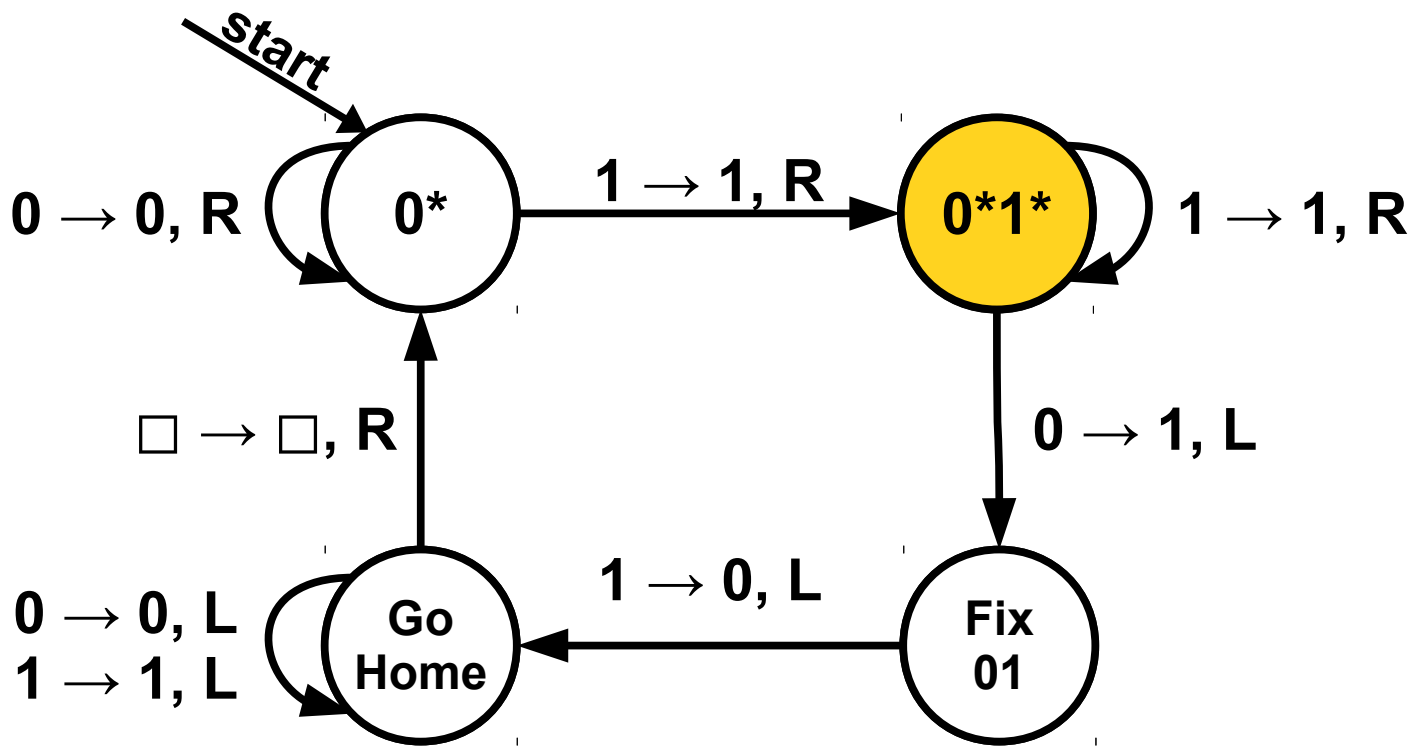


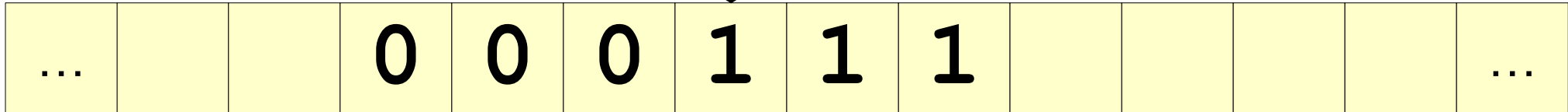
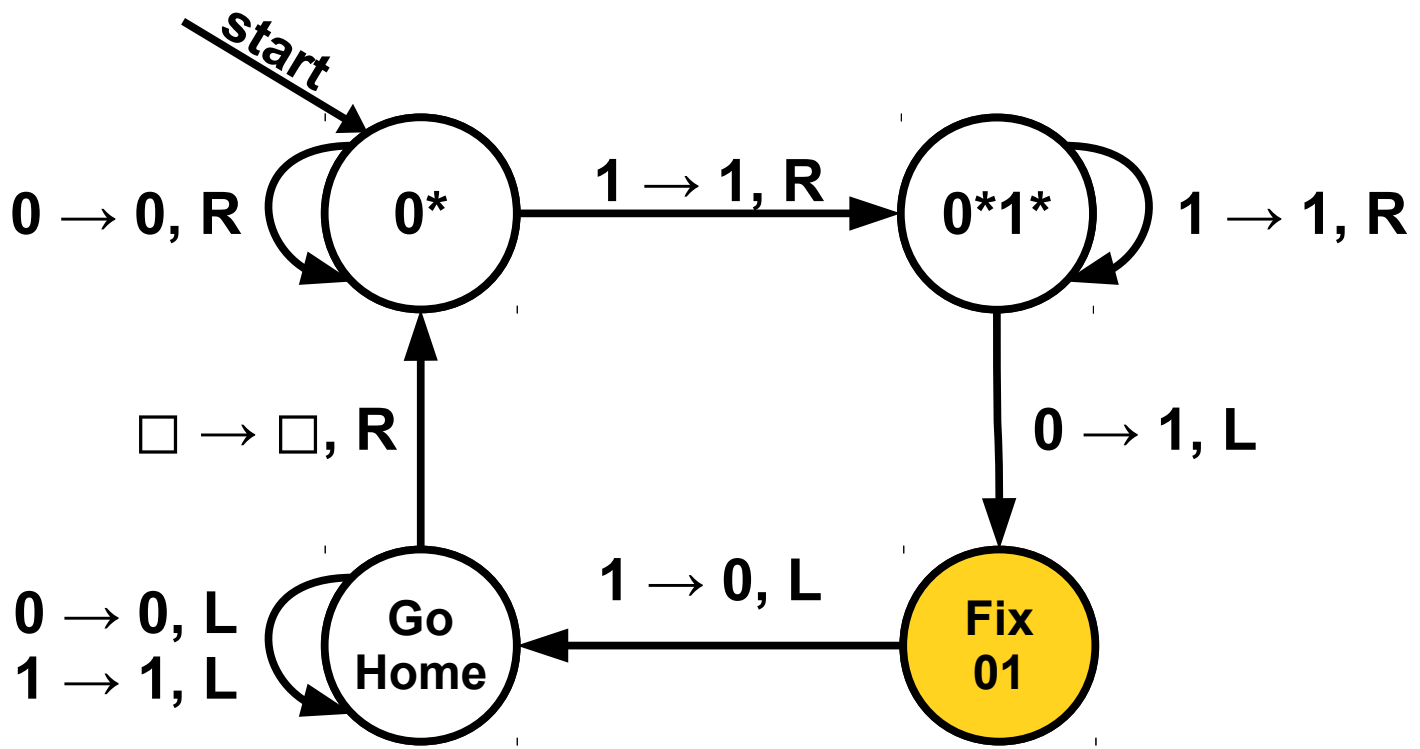




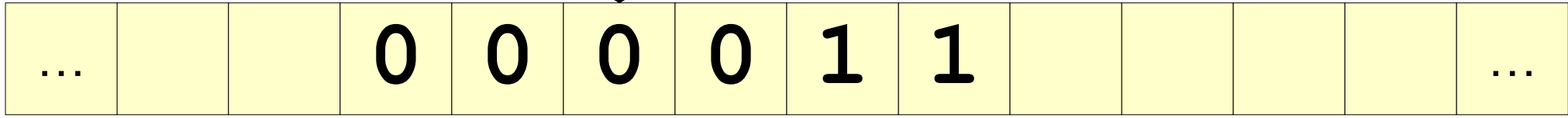
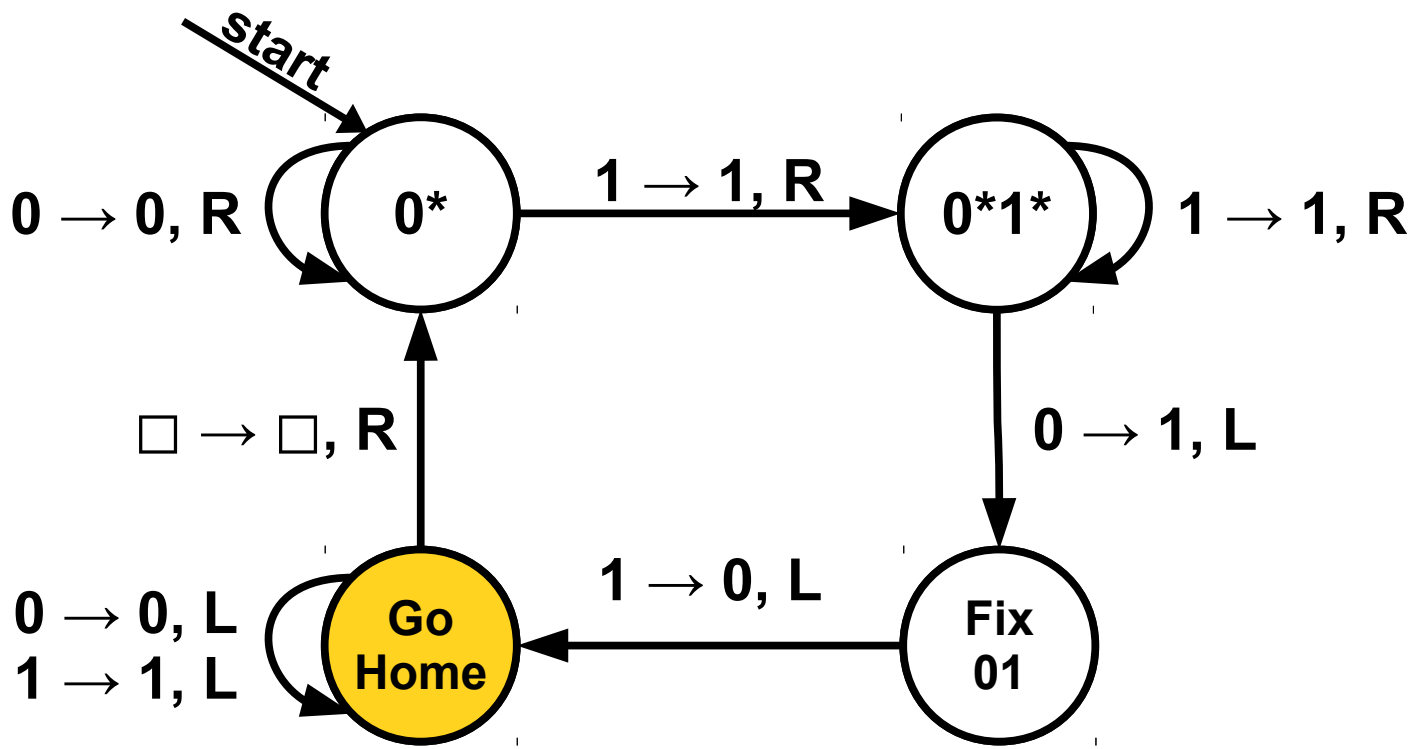


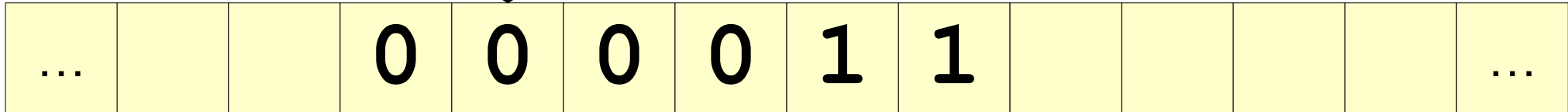
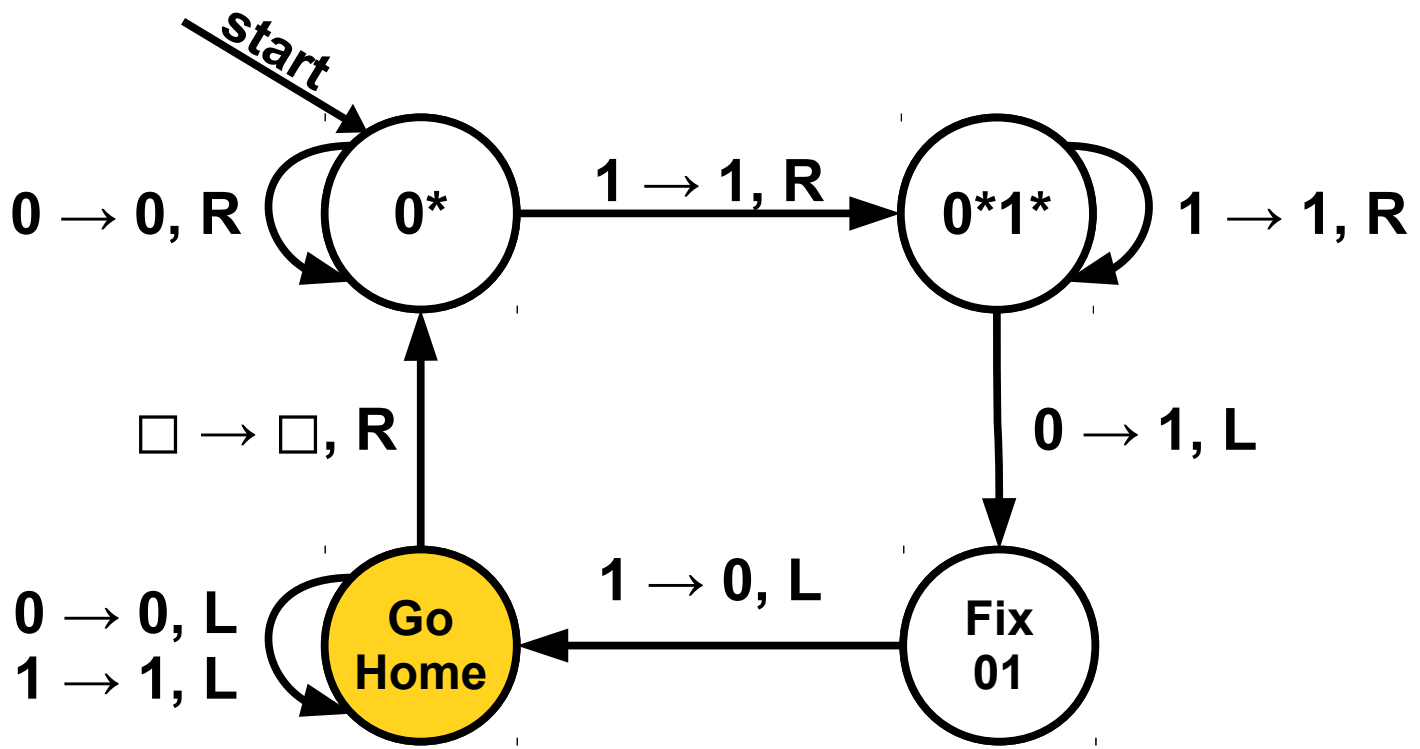


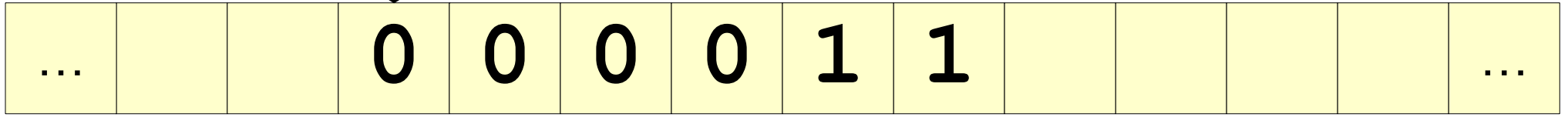
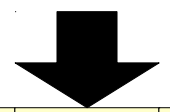
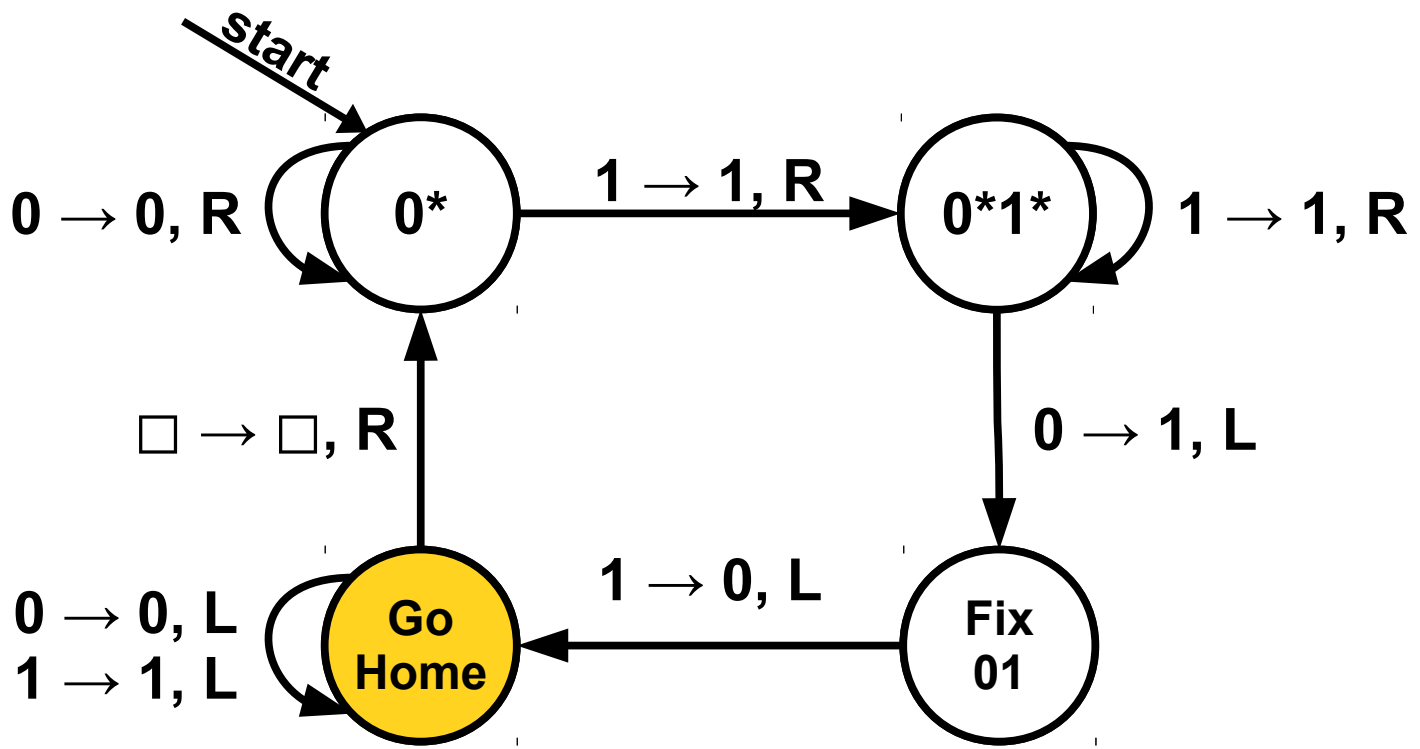


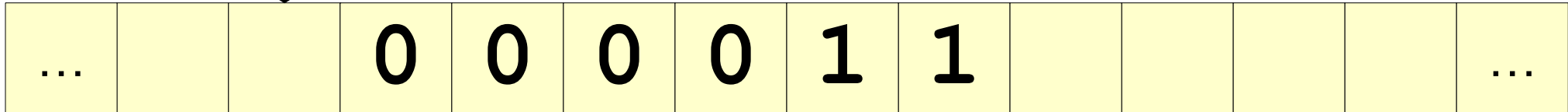
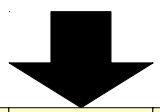
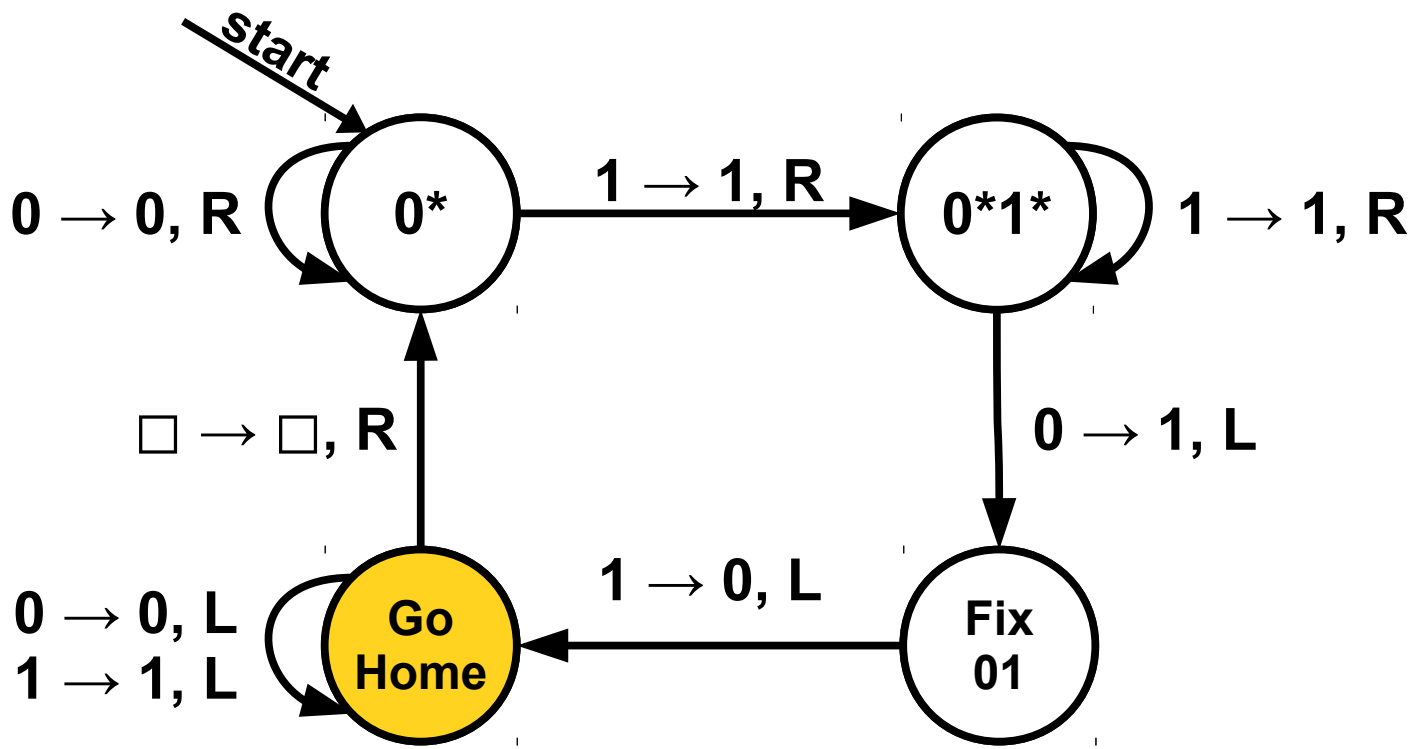


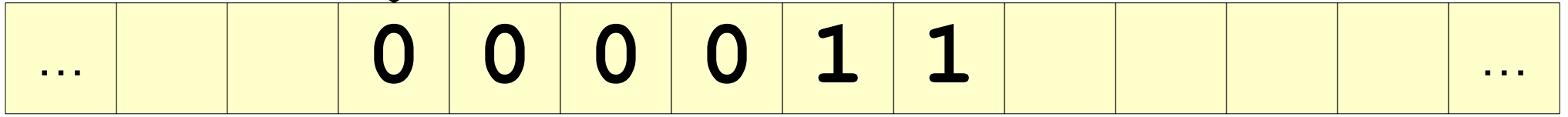
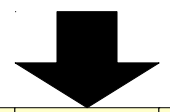
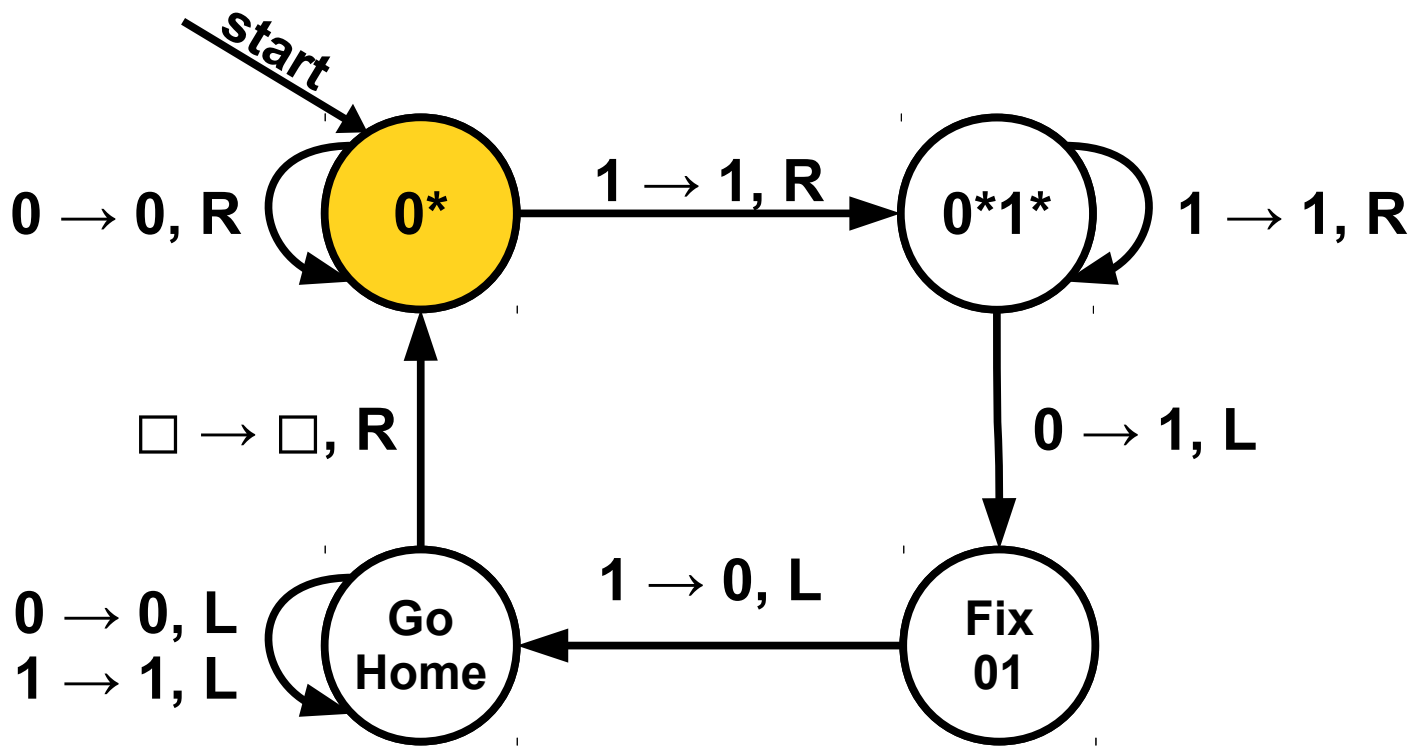


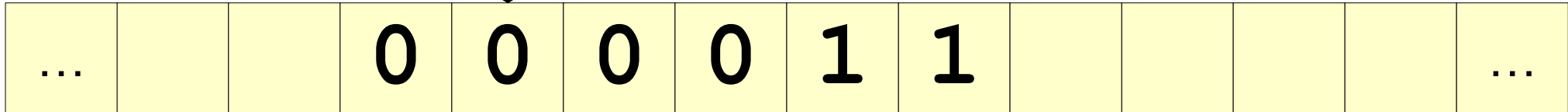
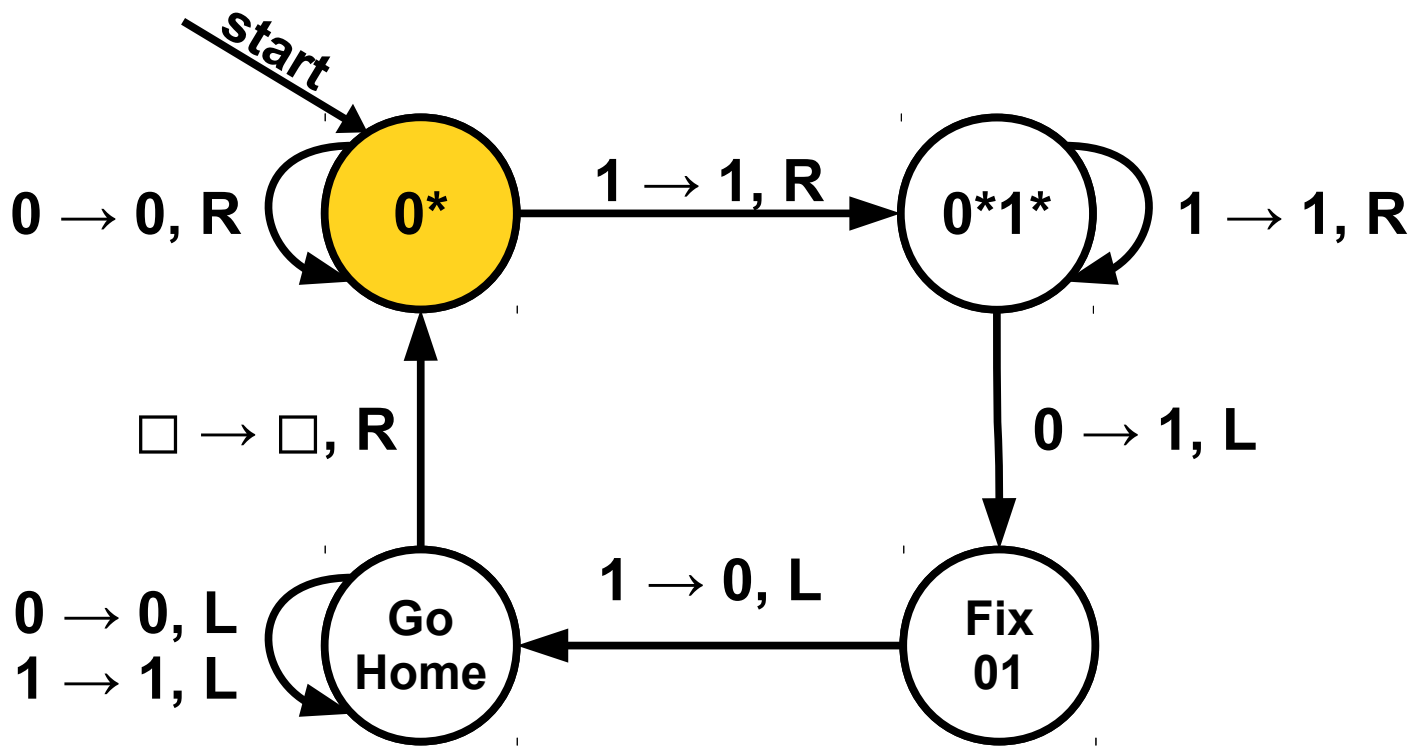


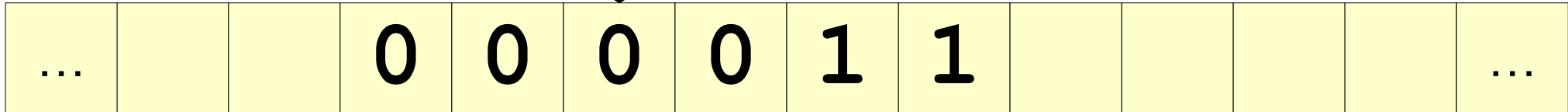
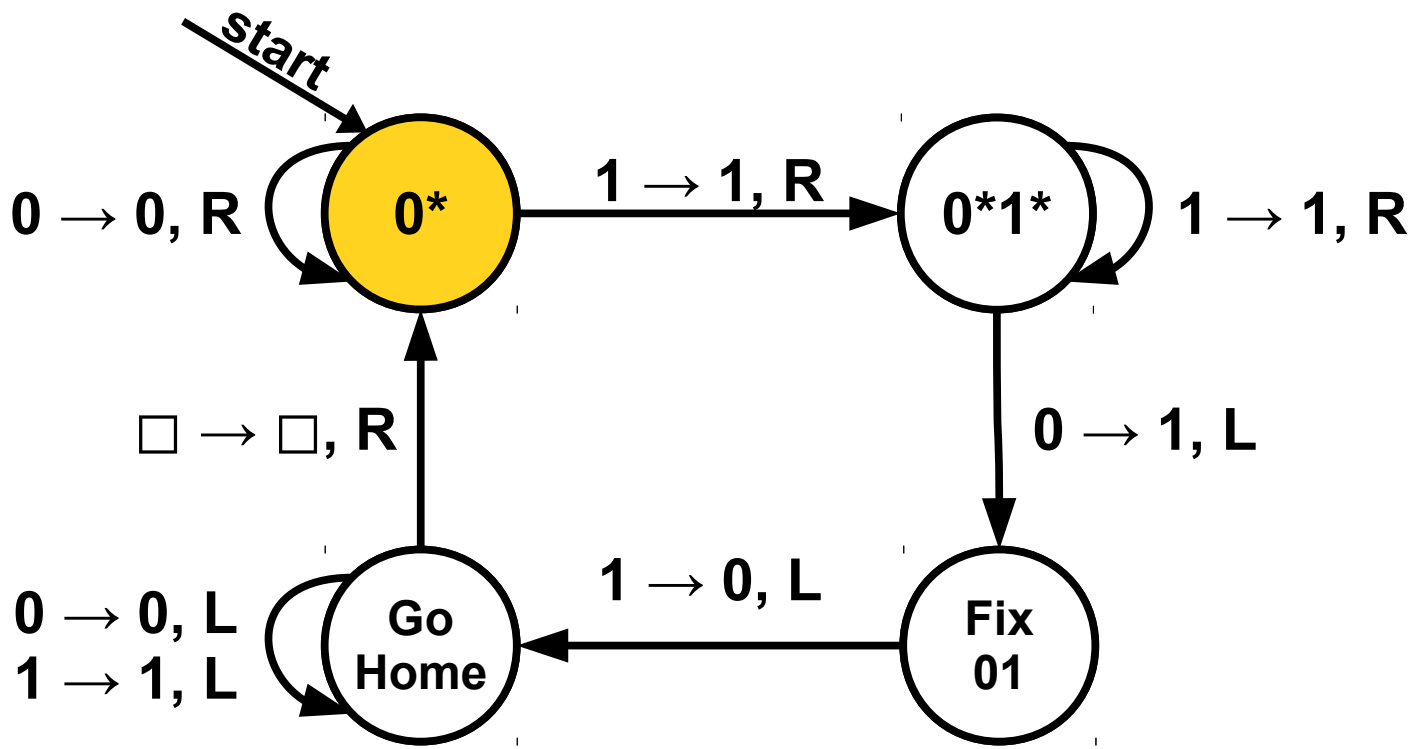


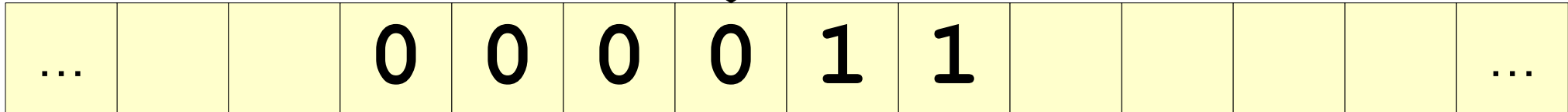
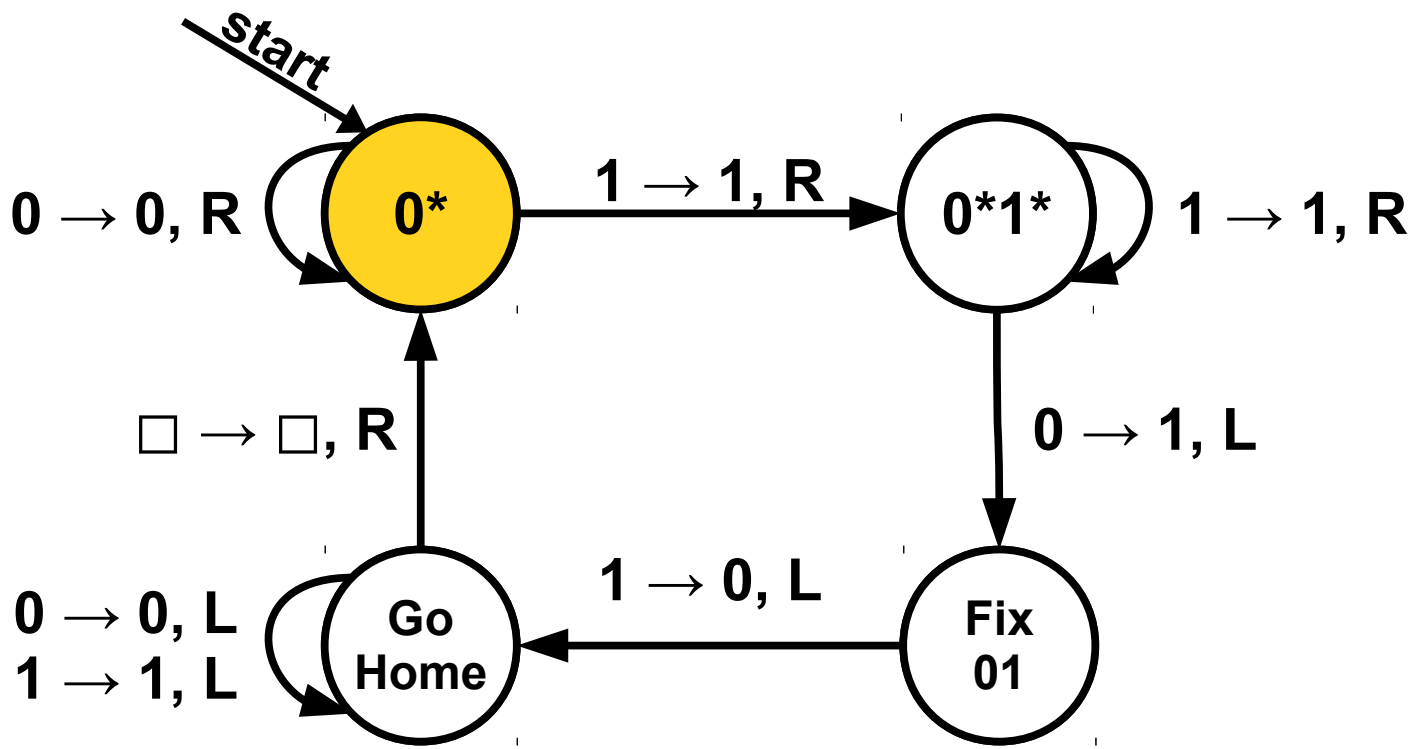




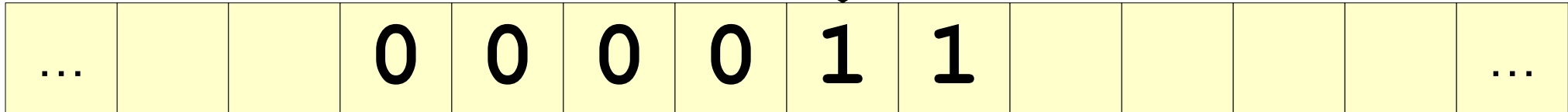
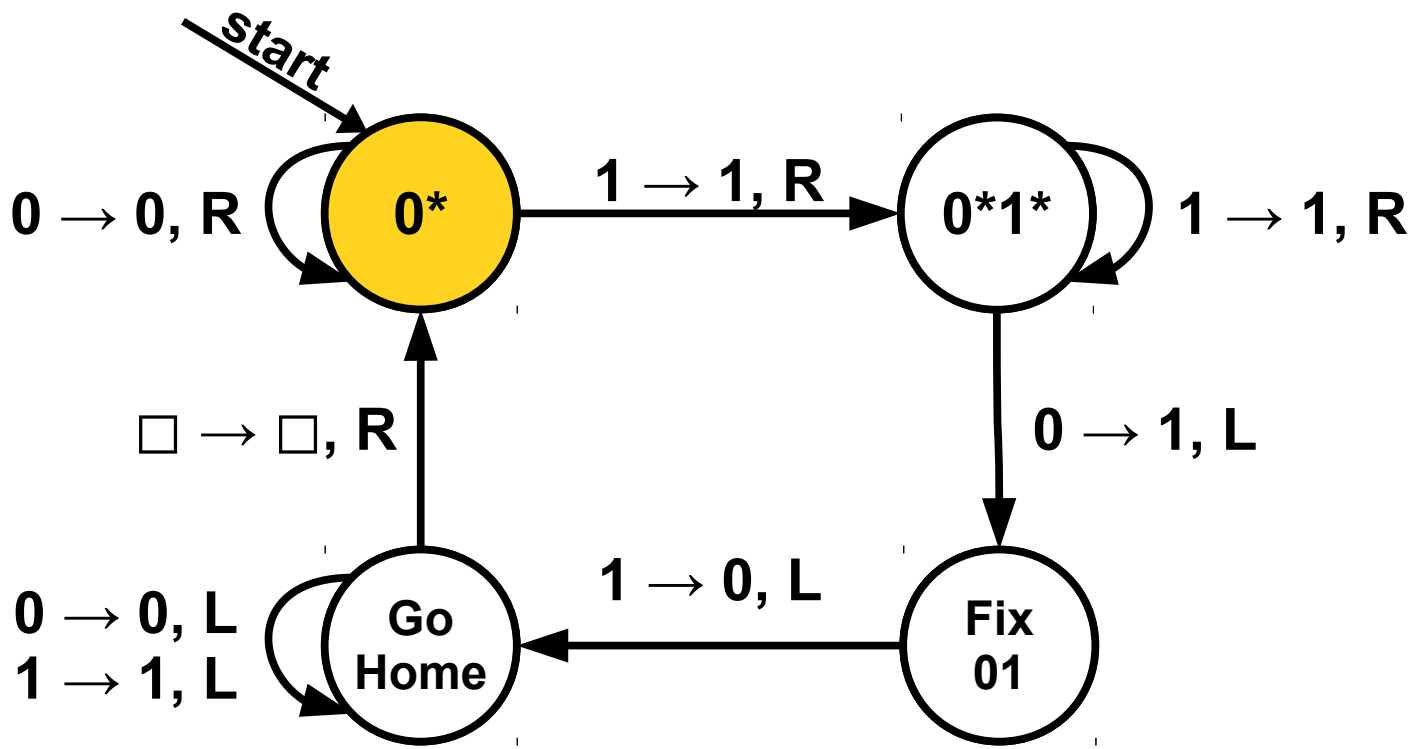


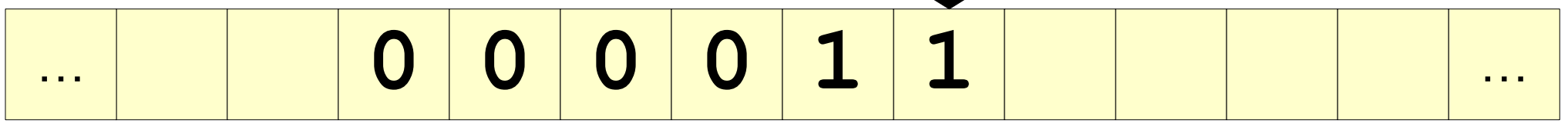
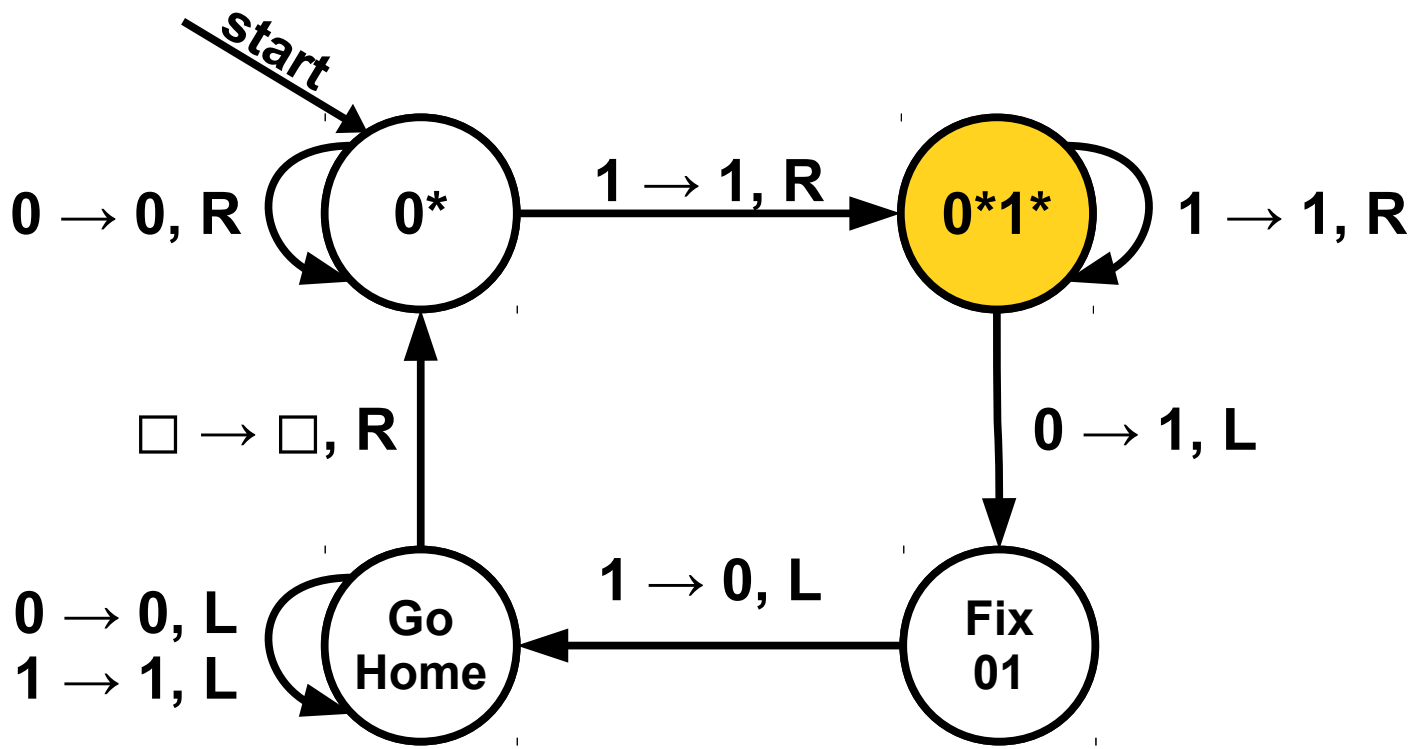


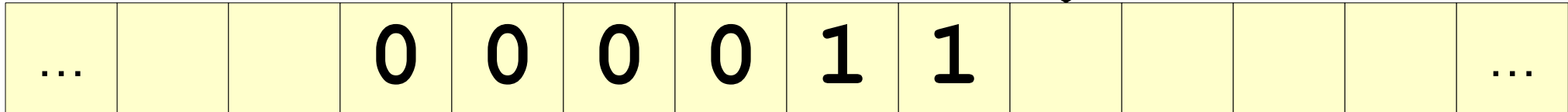
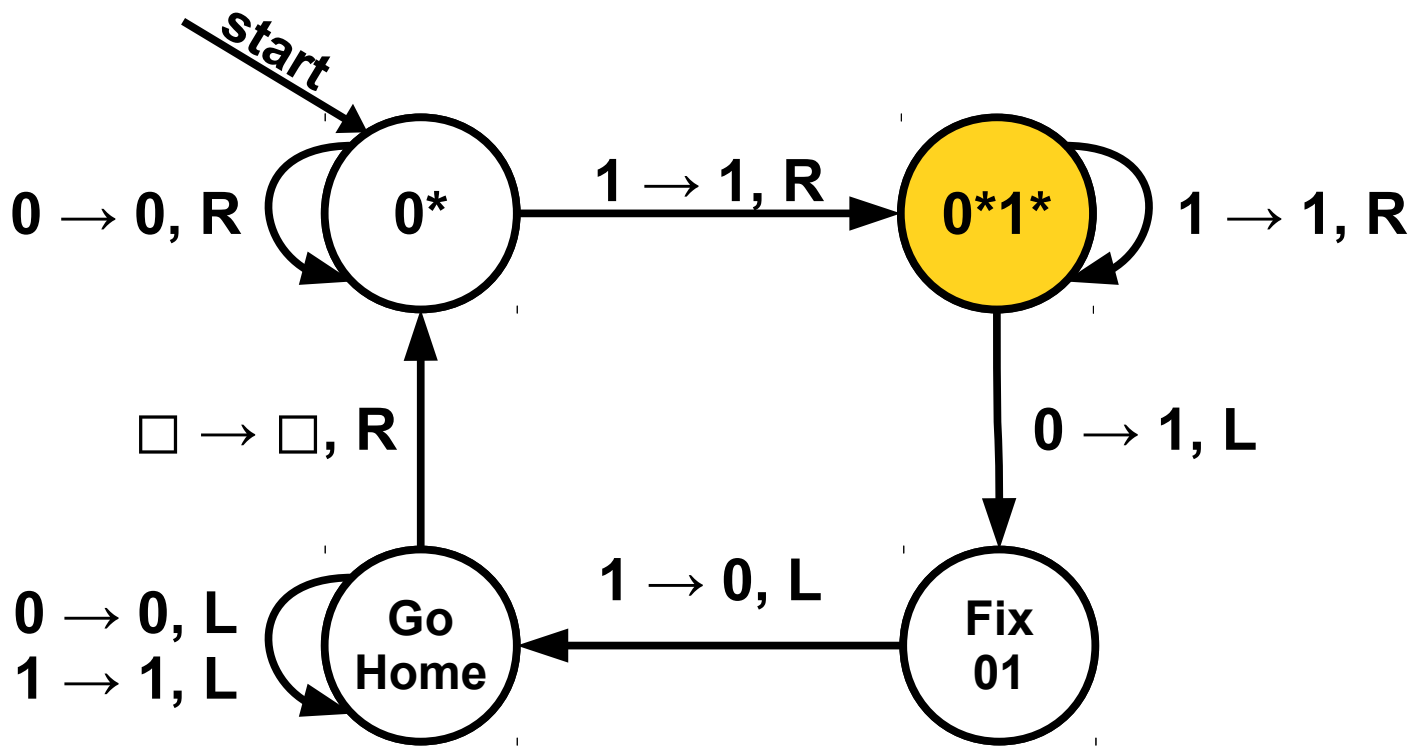




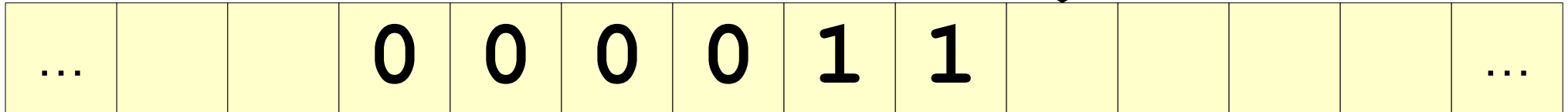
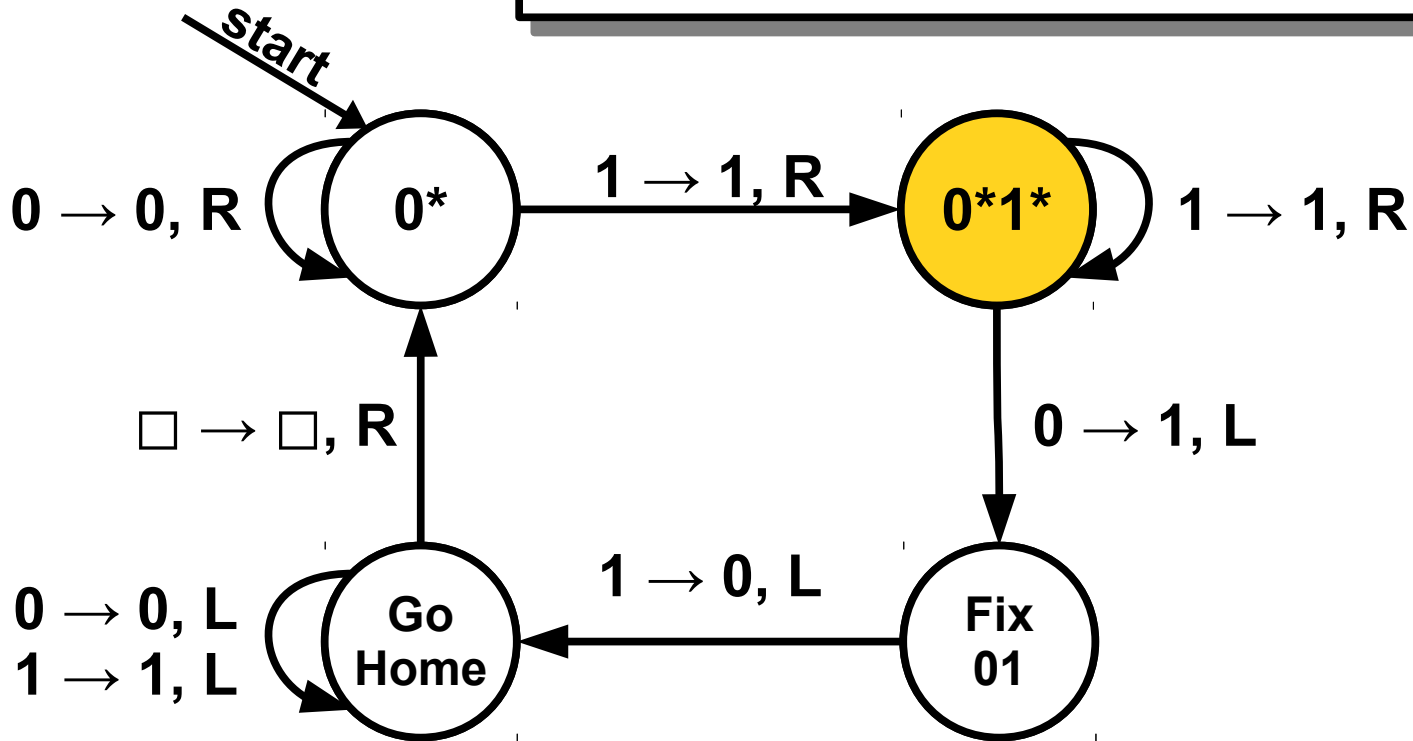


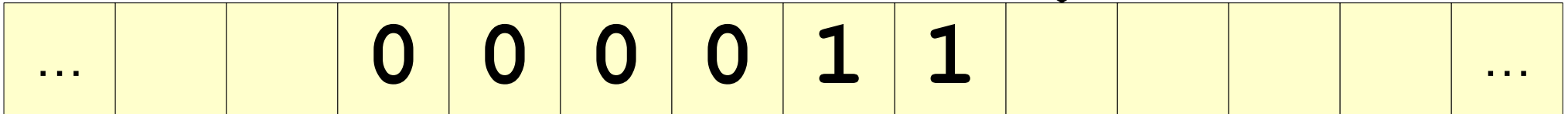
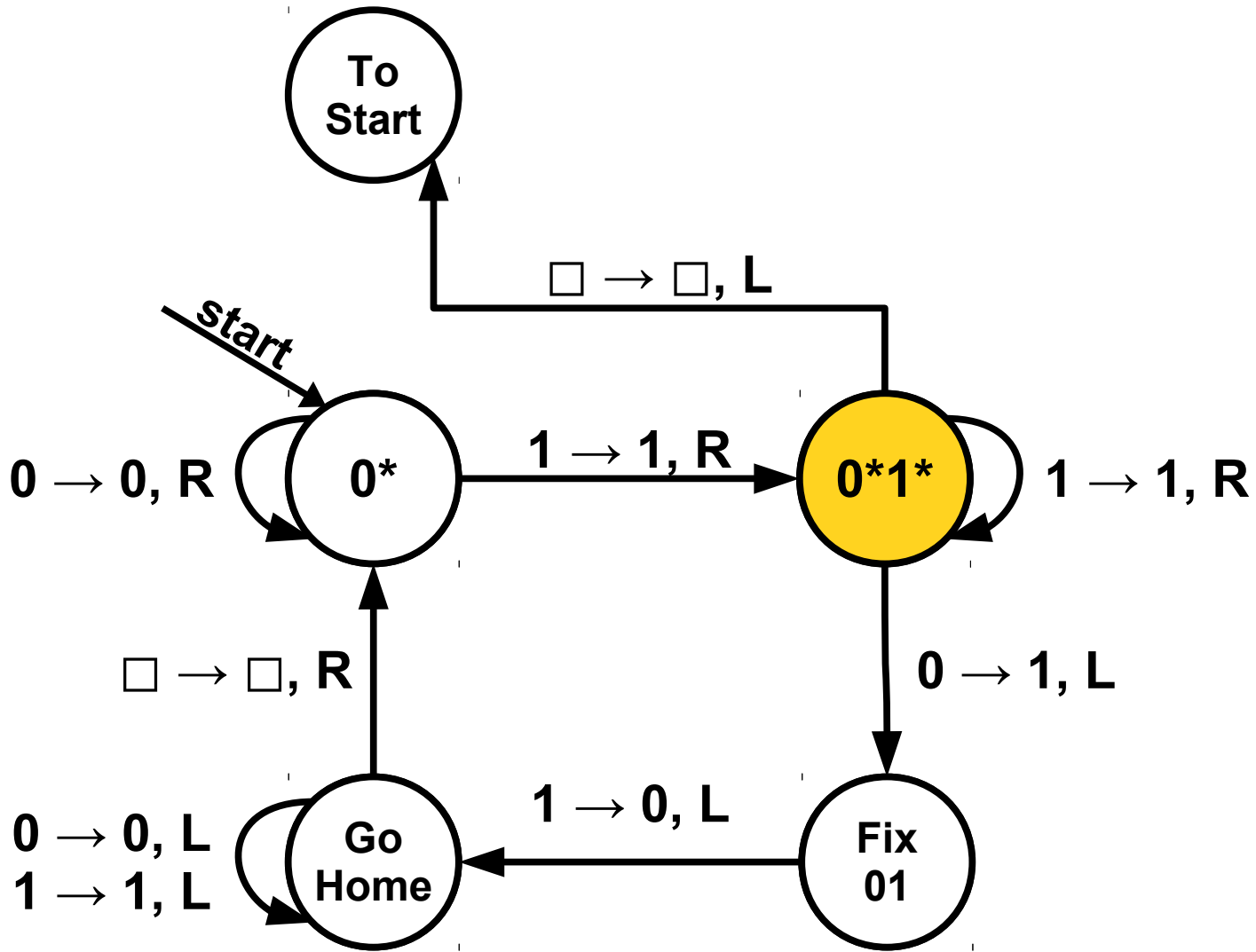


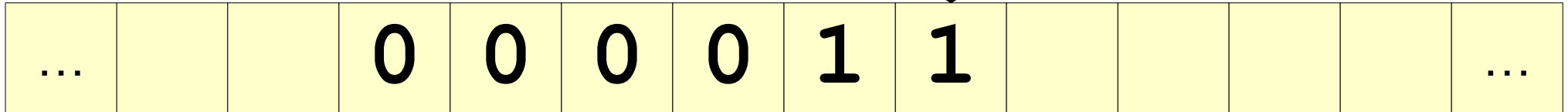
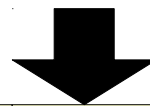
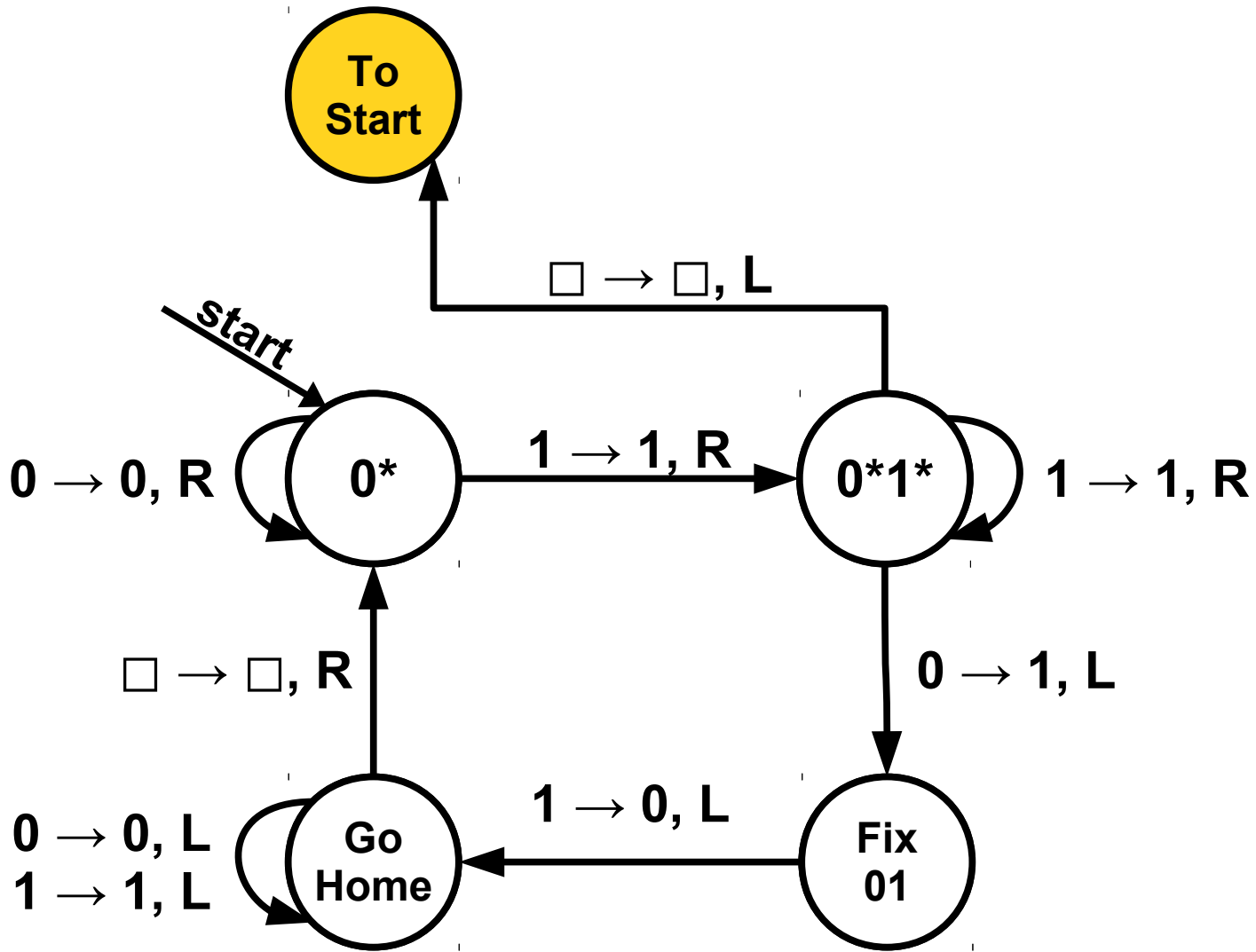


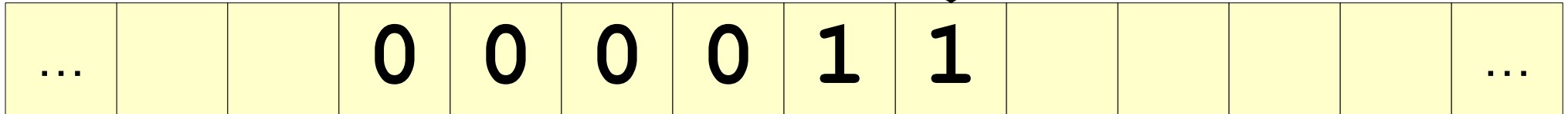
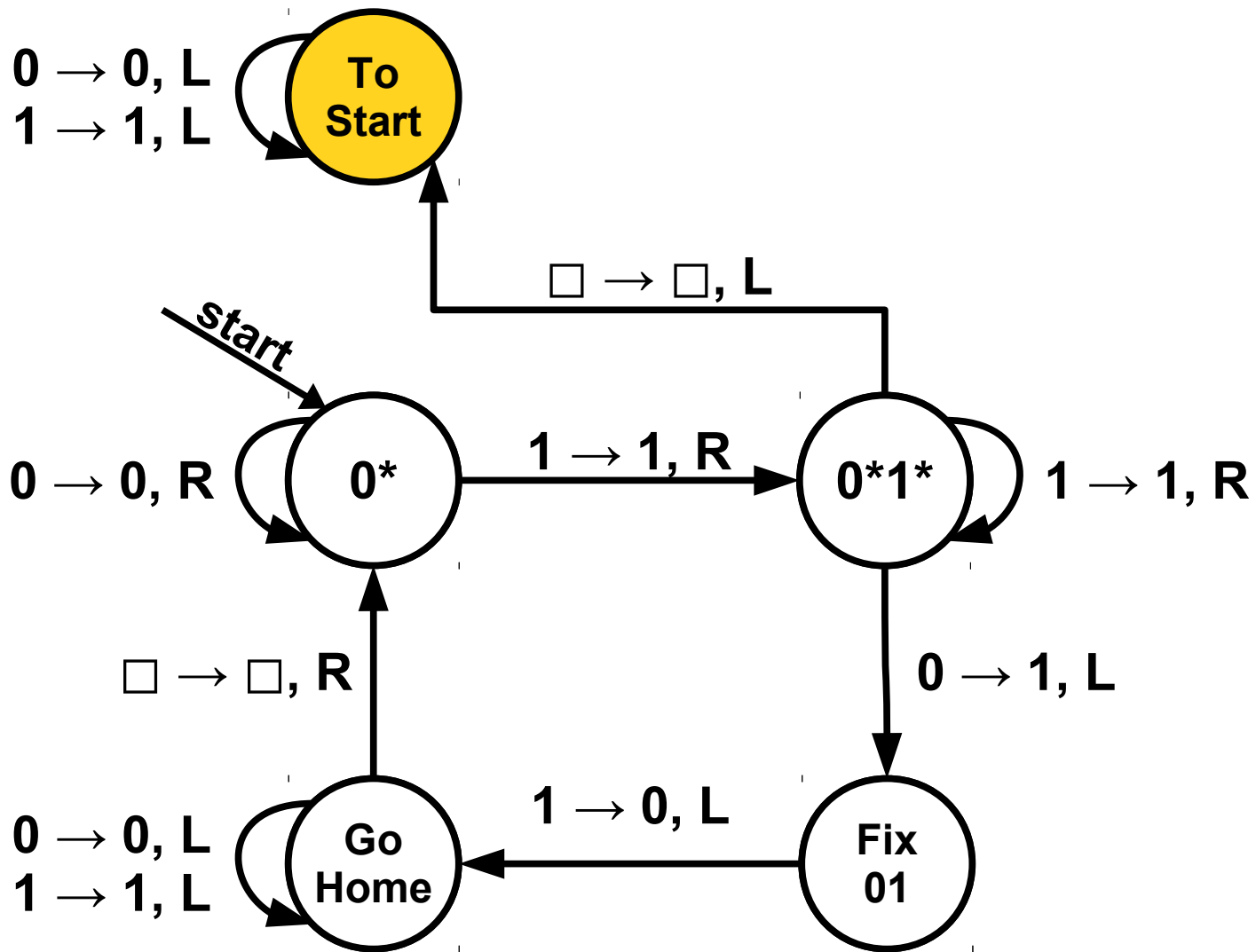


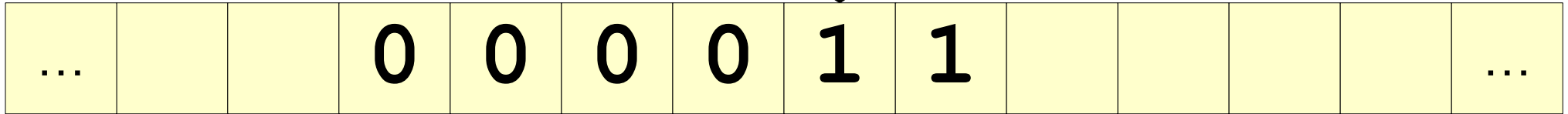
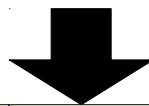
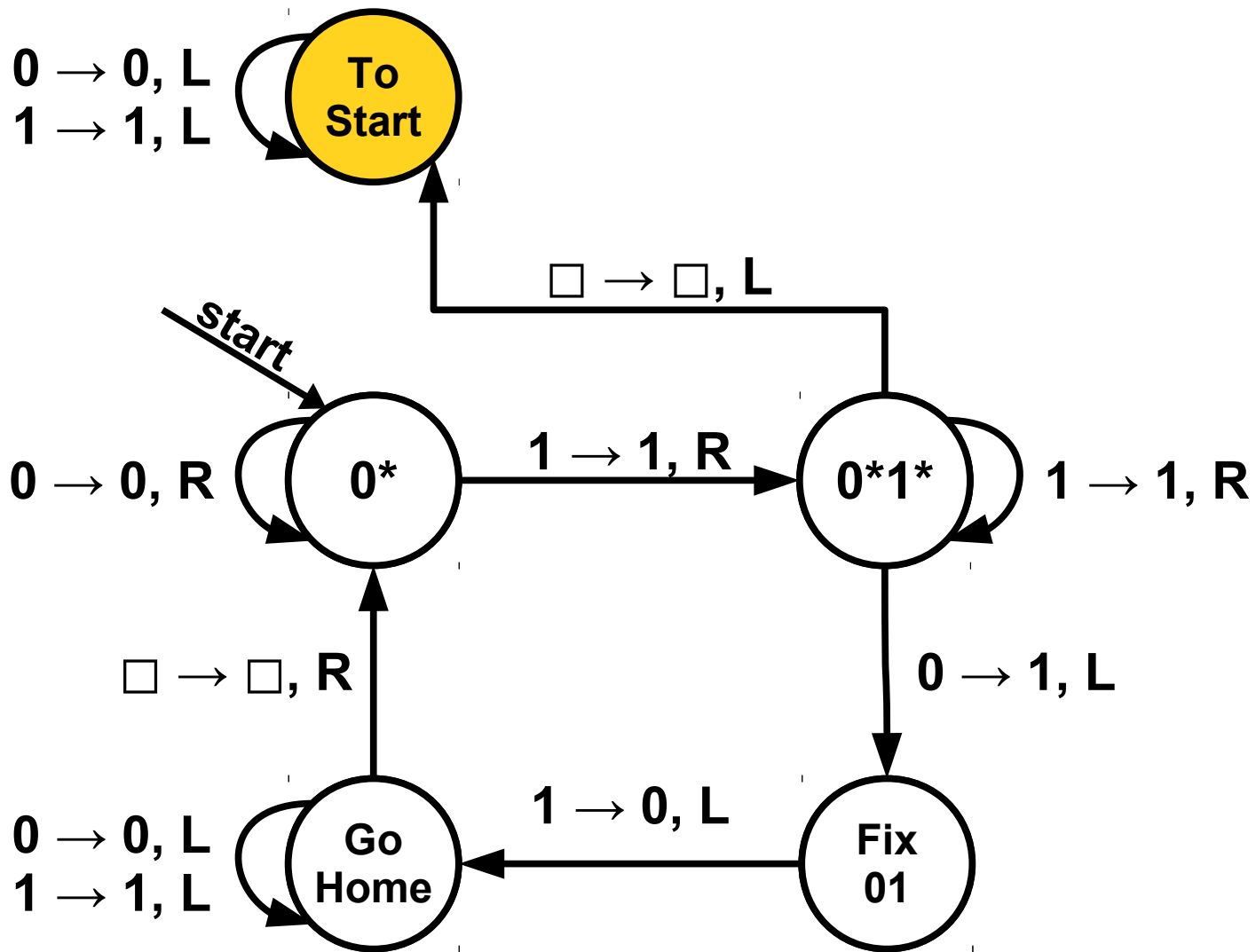
Our ultimate goal here was to sort everything so we could hand it off to the machine to check for  $0^n 1^n$ . Let's rewind the tape head back to the start.



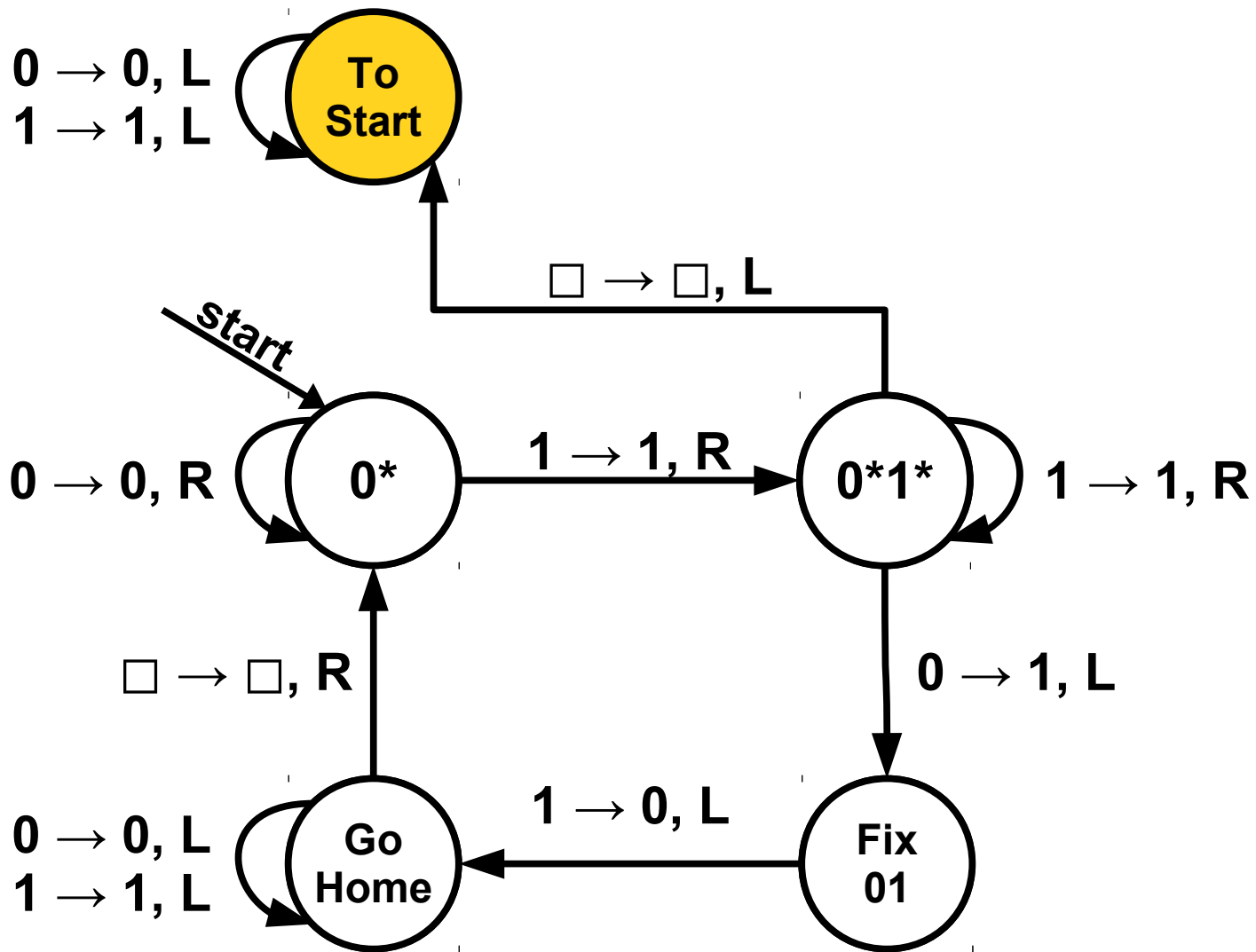


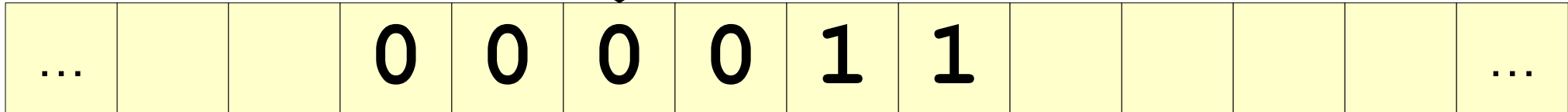
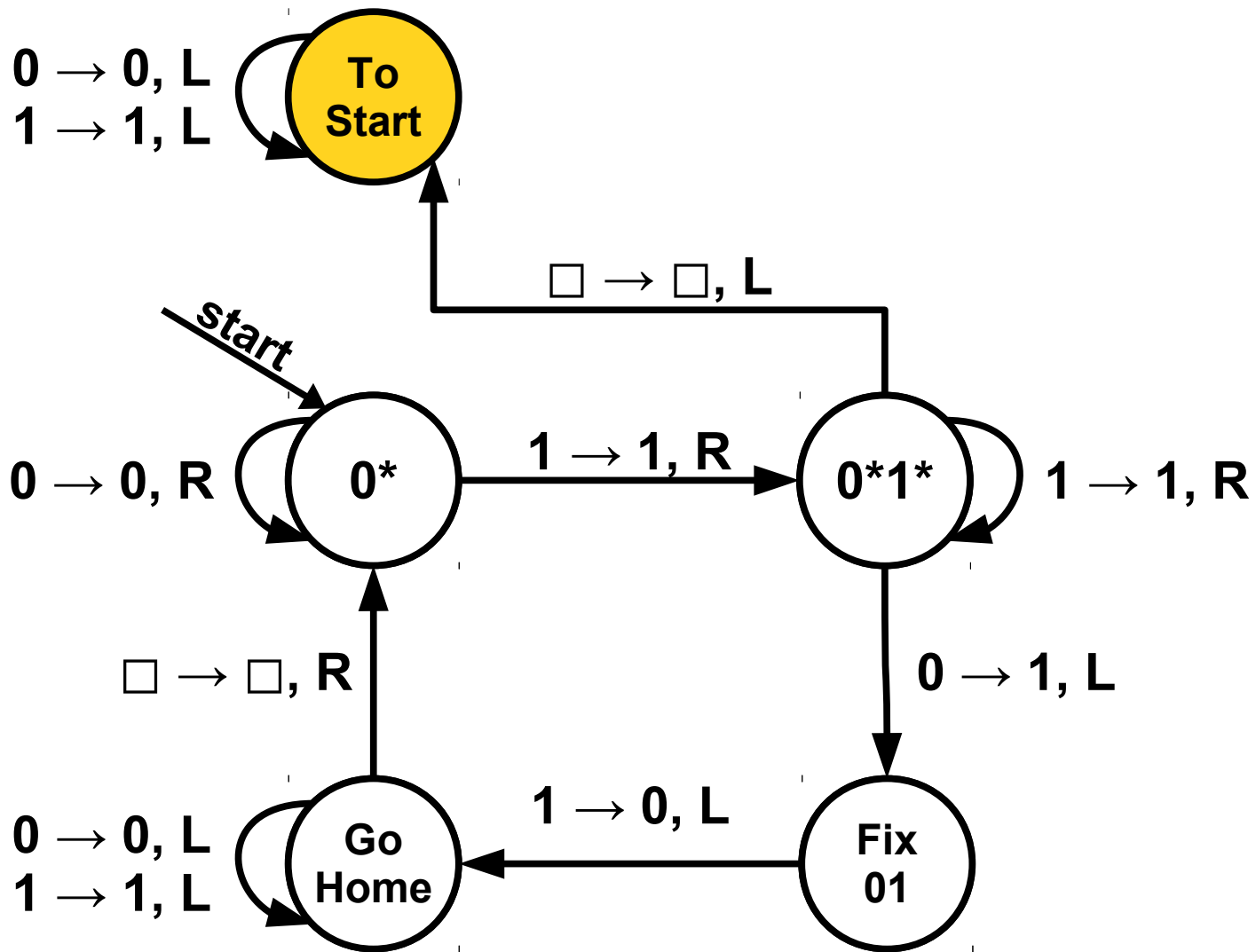


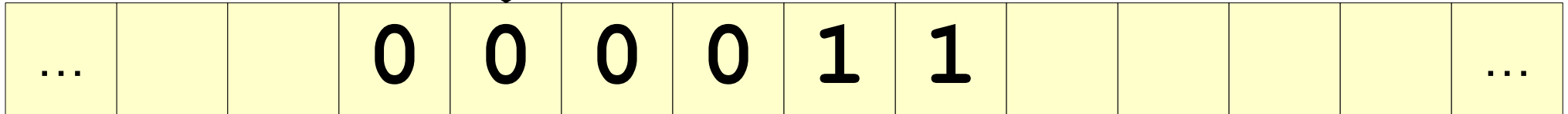
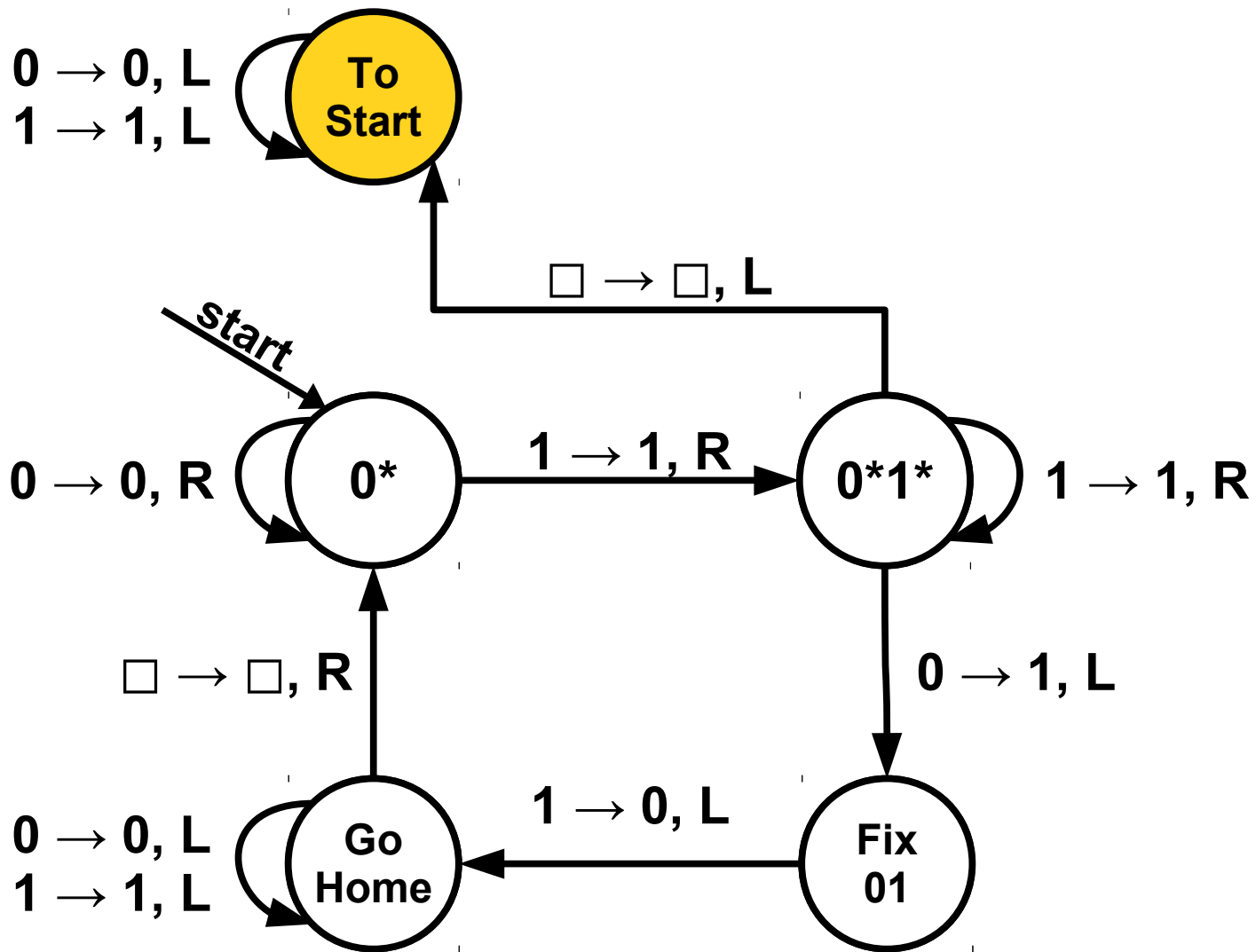


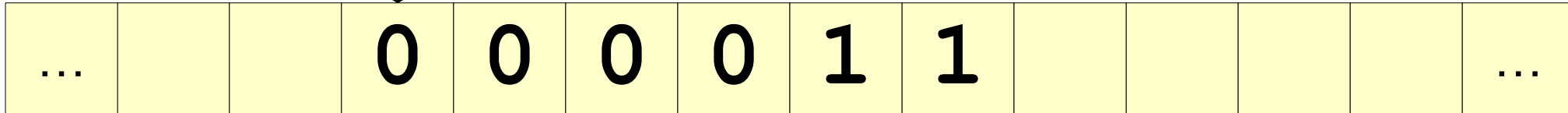
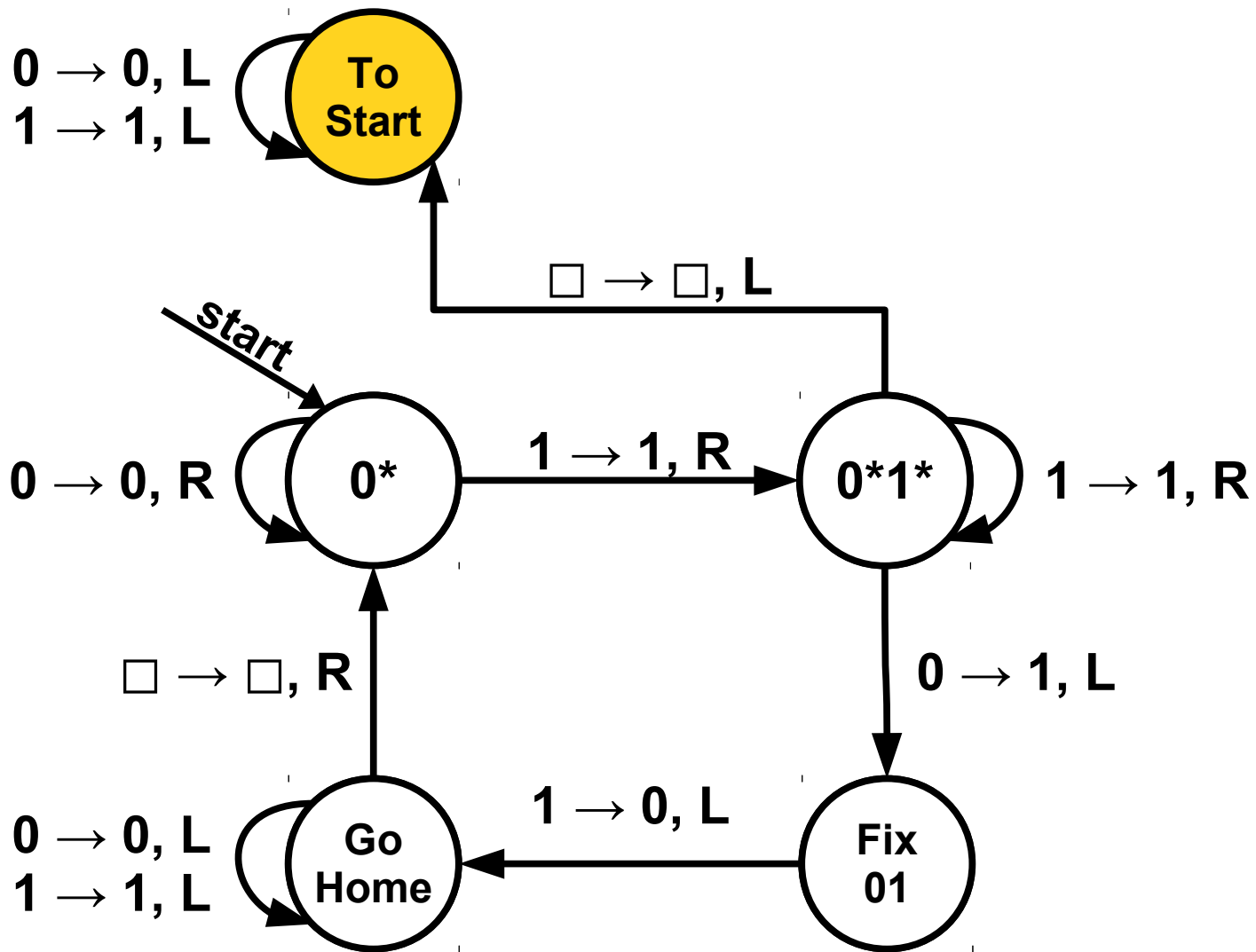


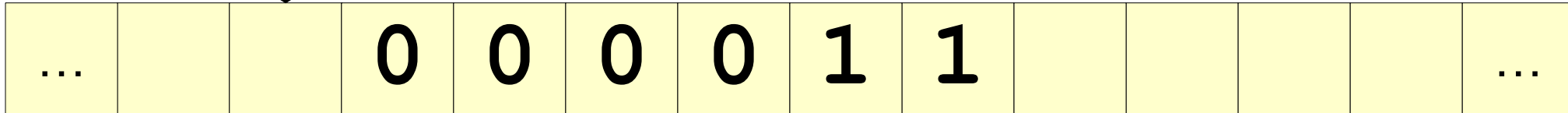
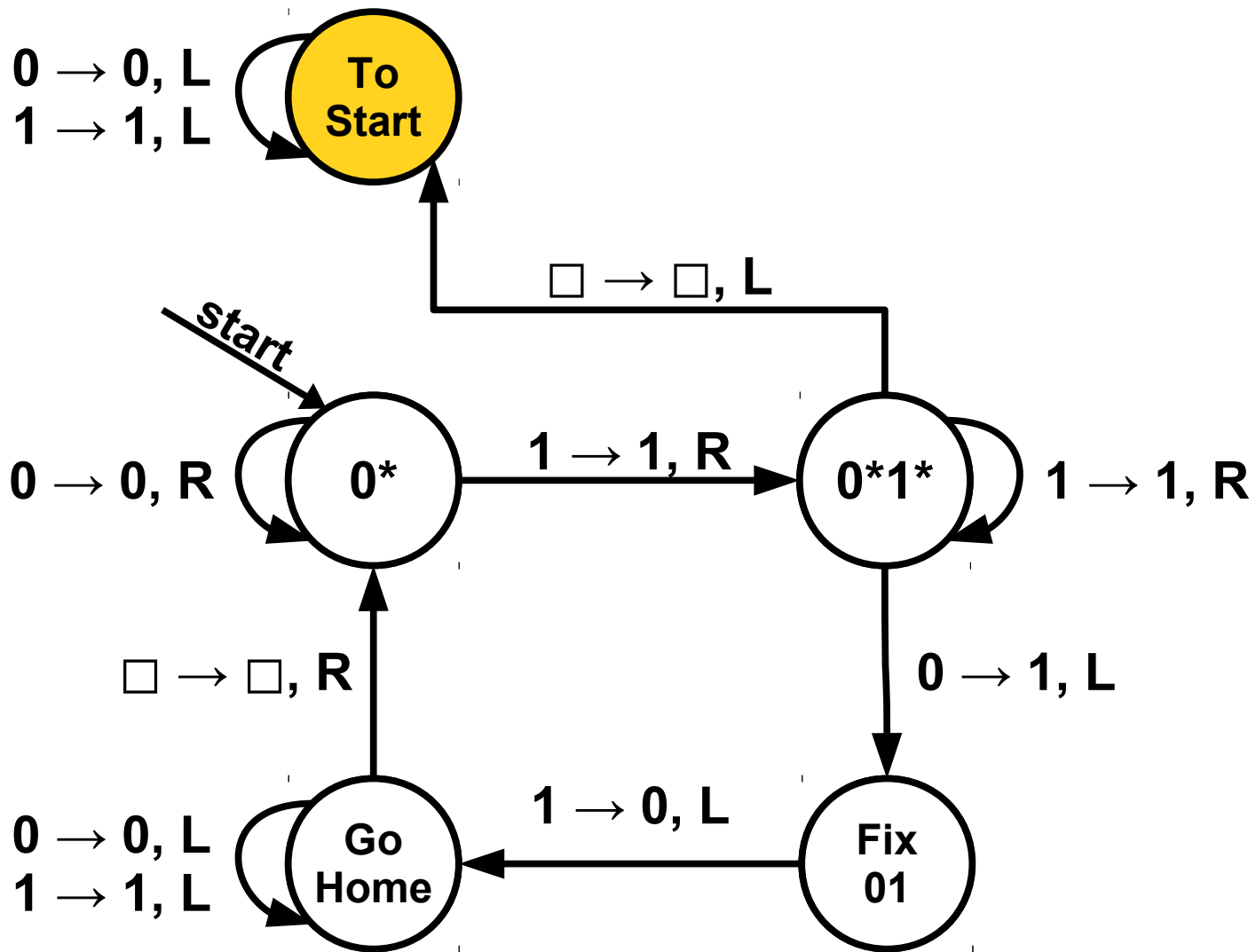


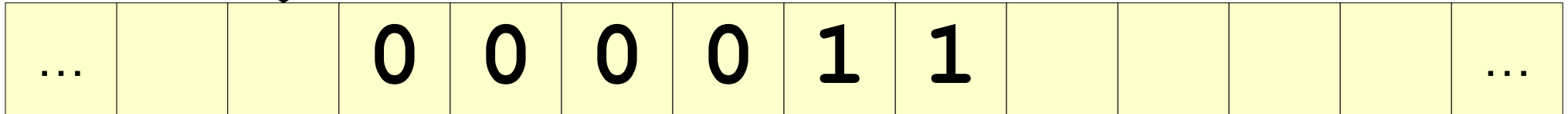
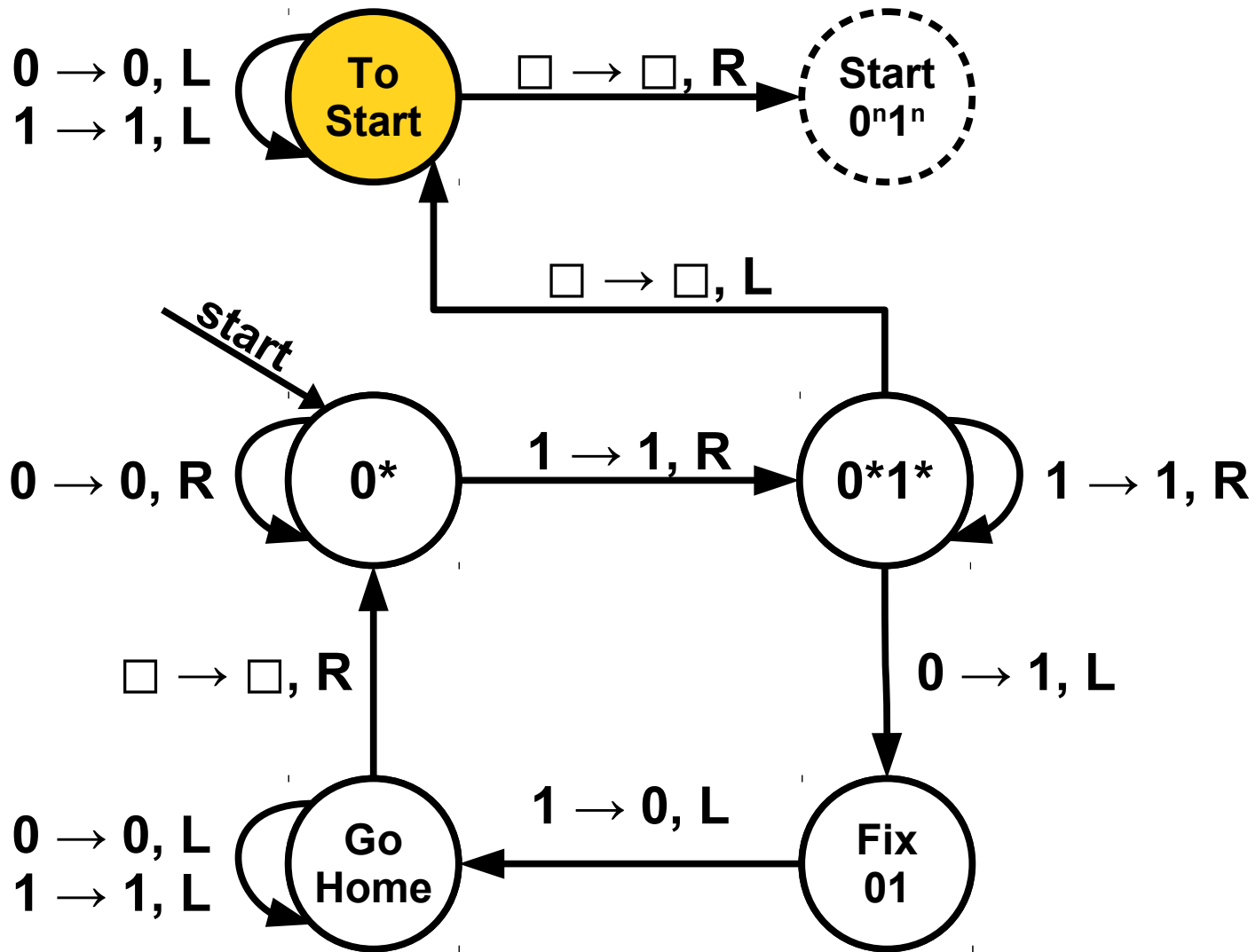


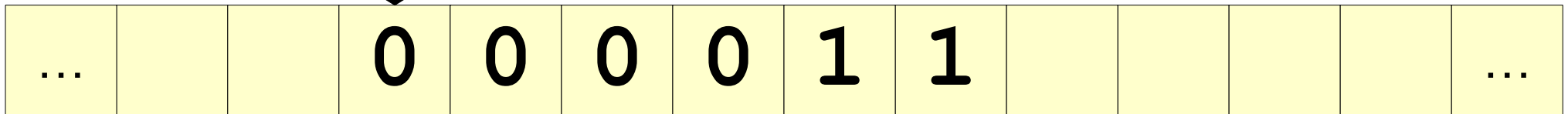
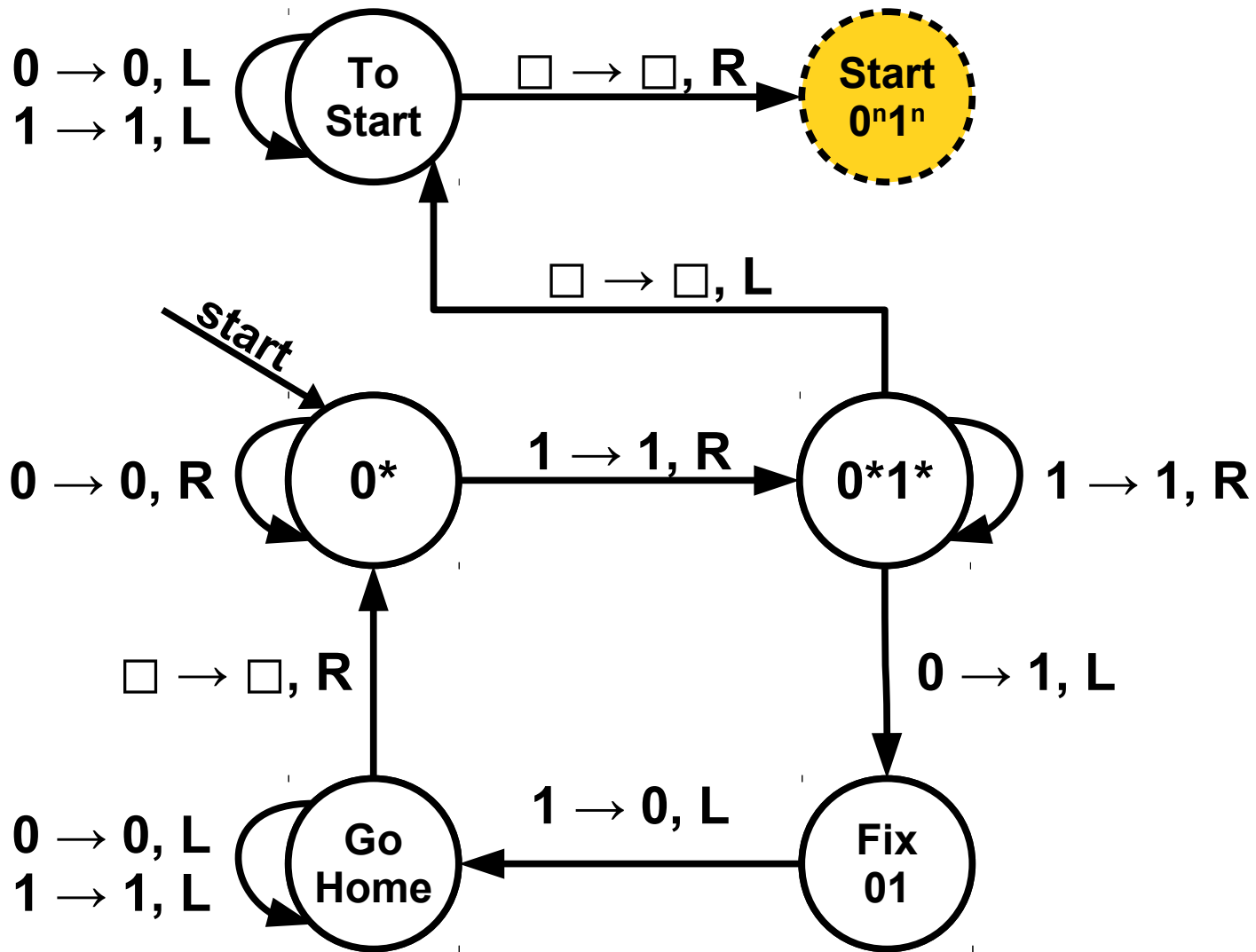


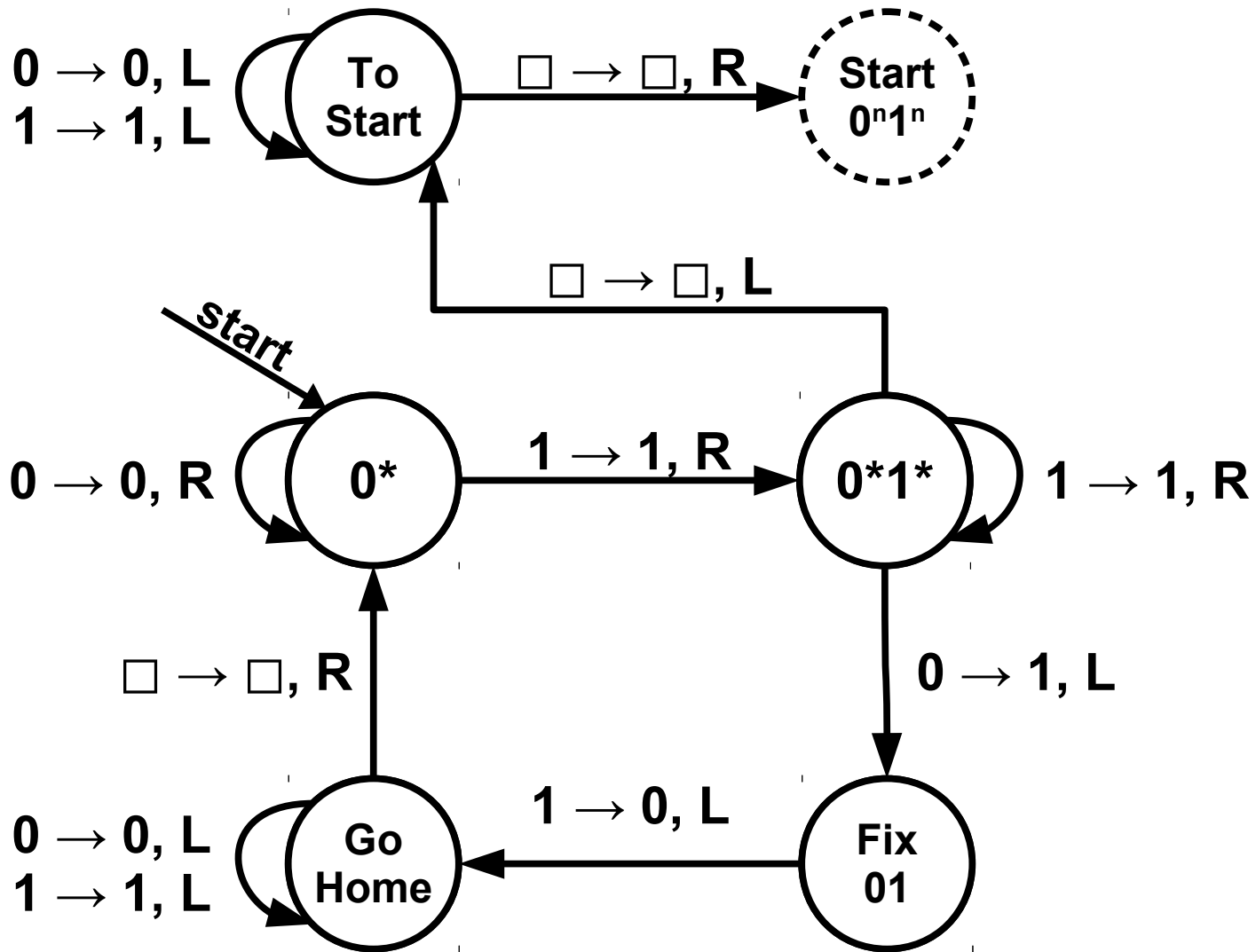




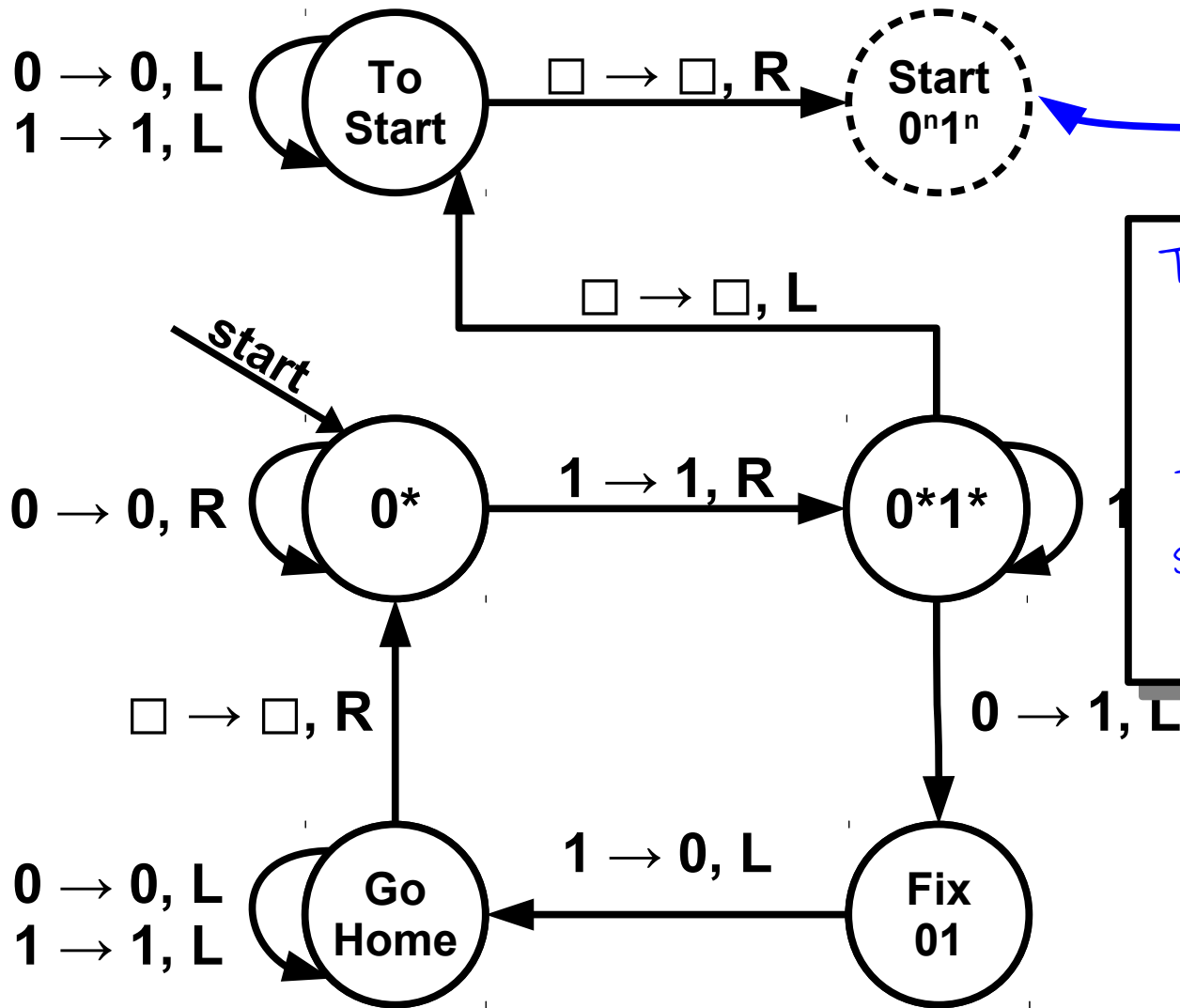




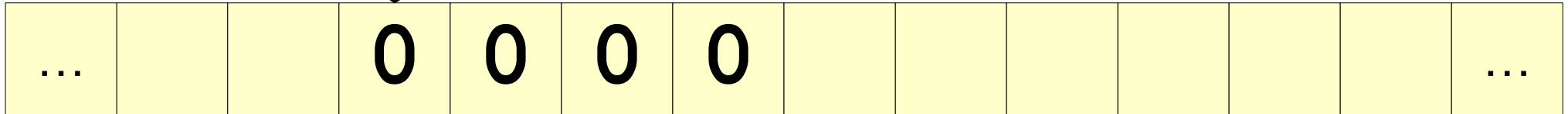
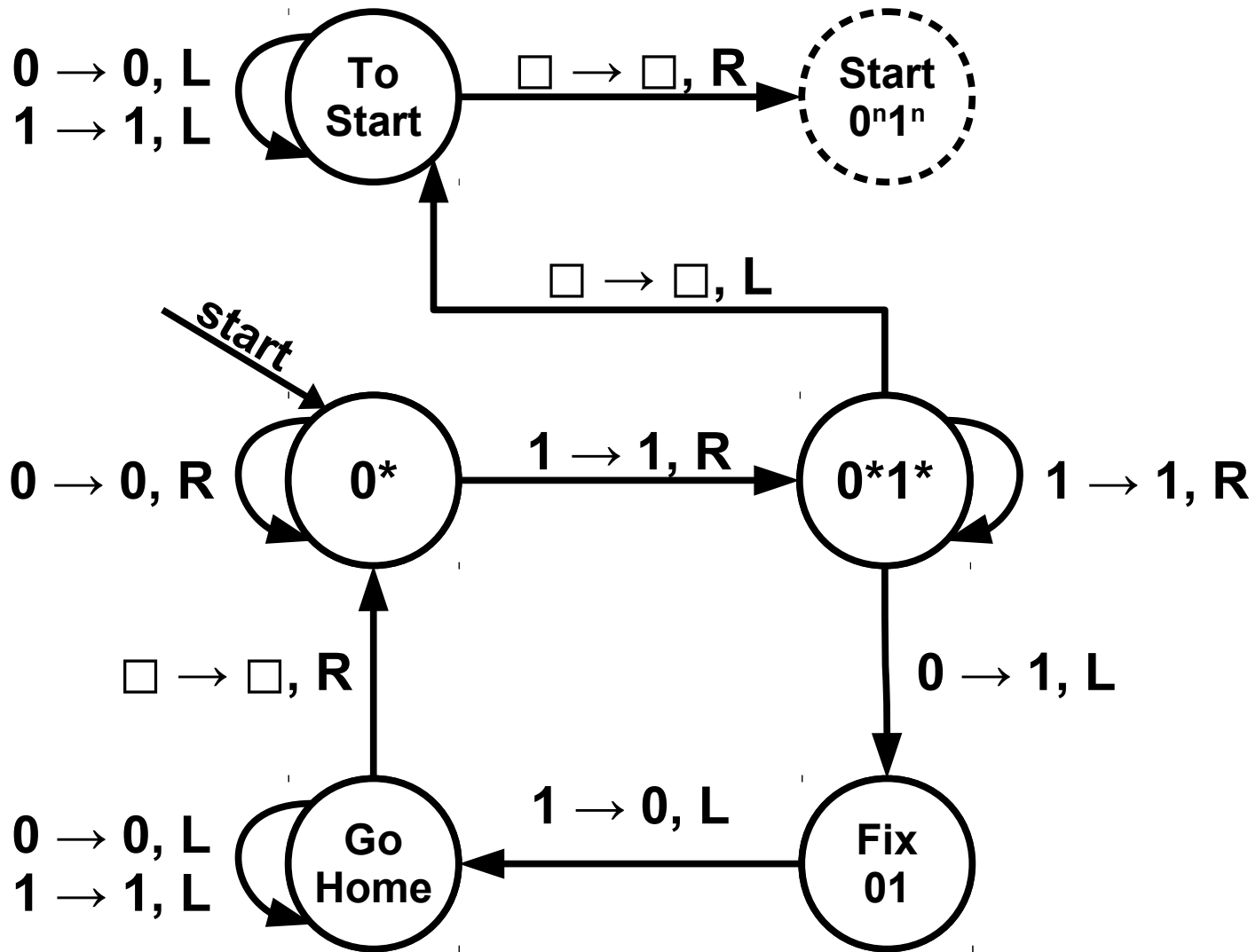


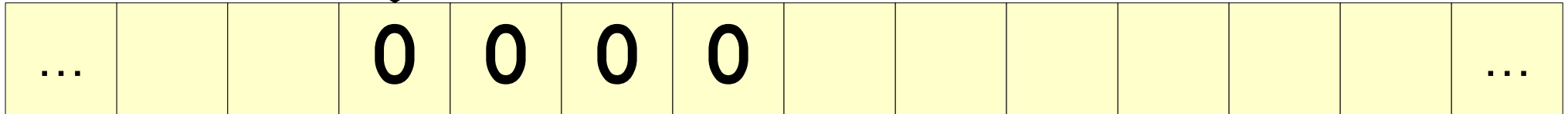
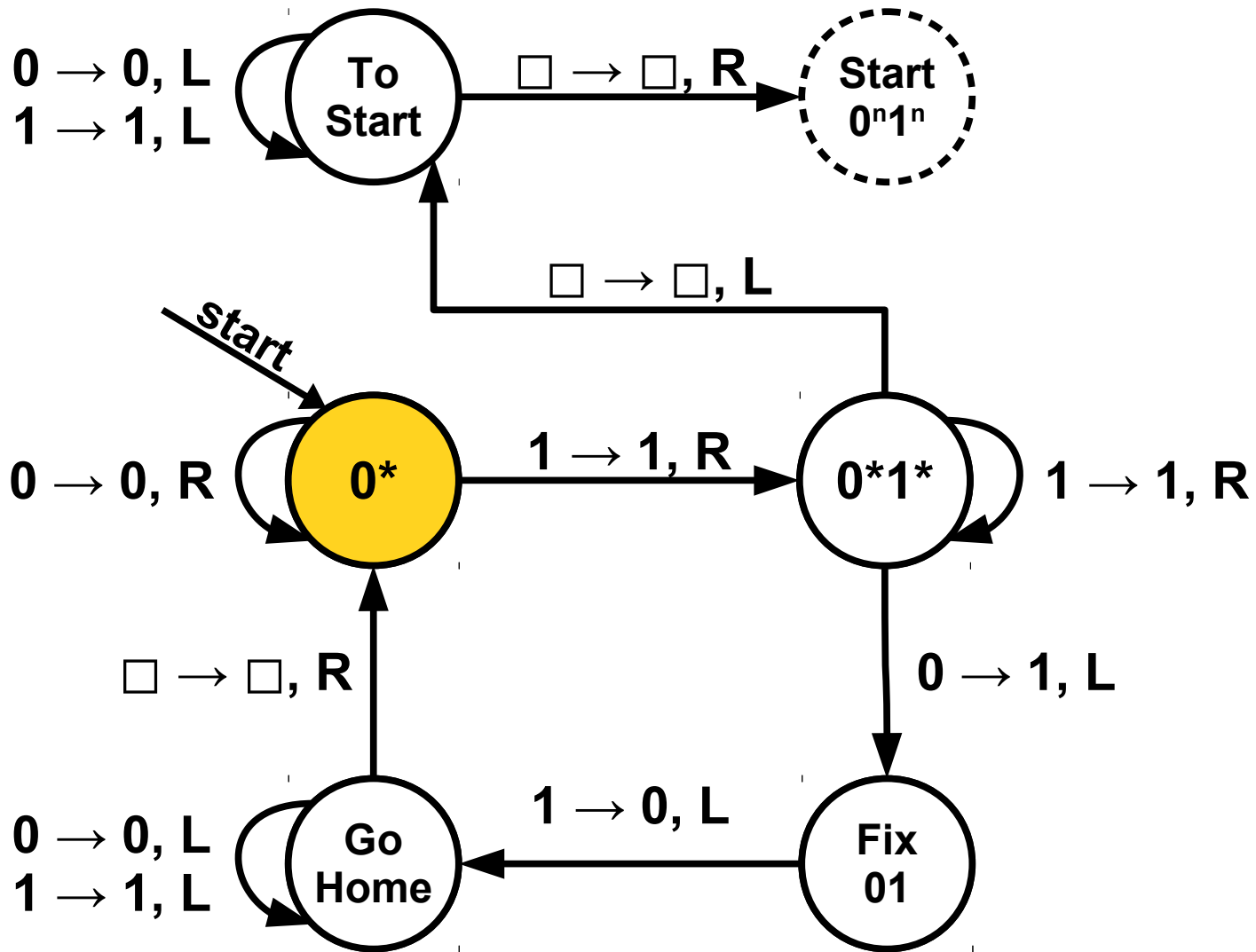


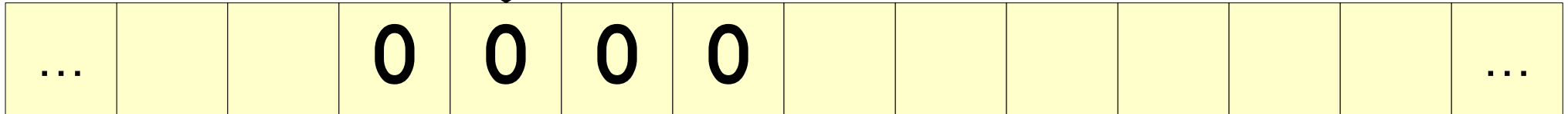
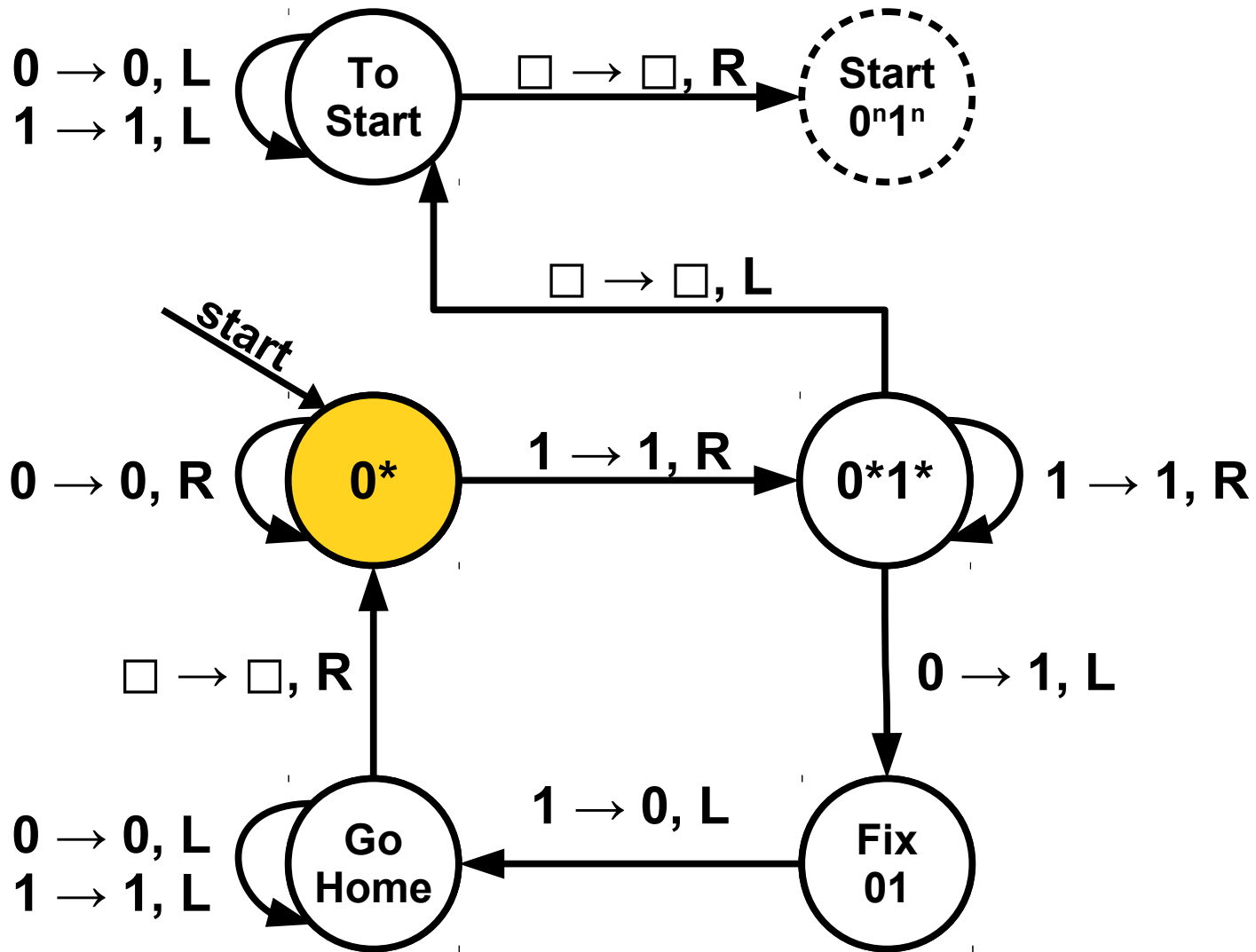




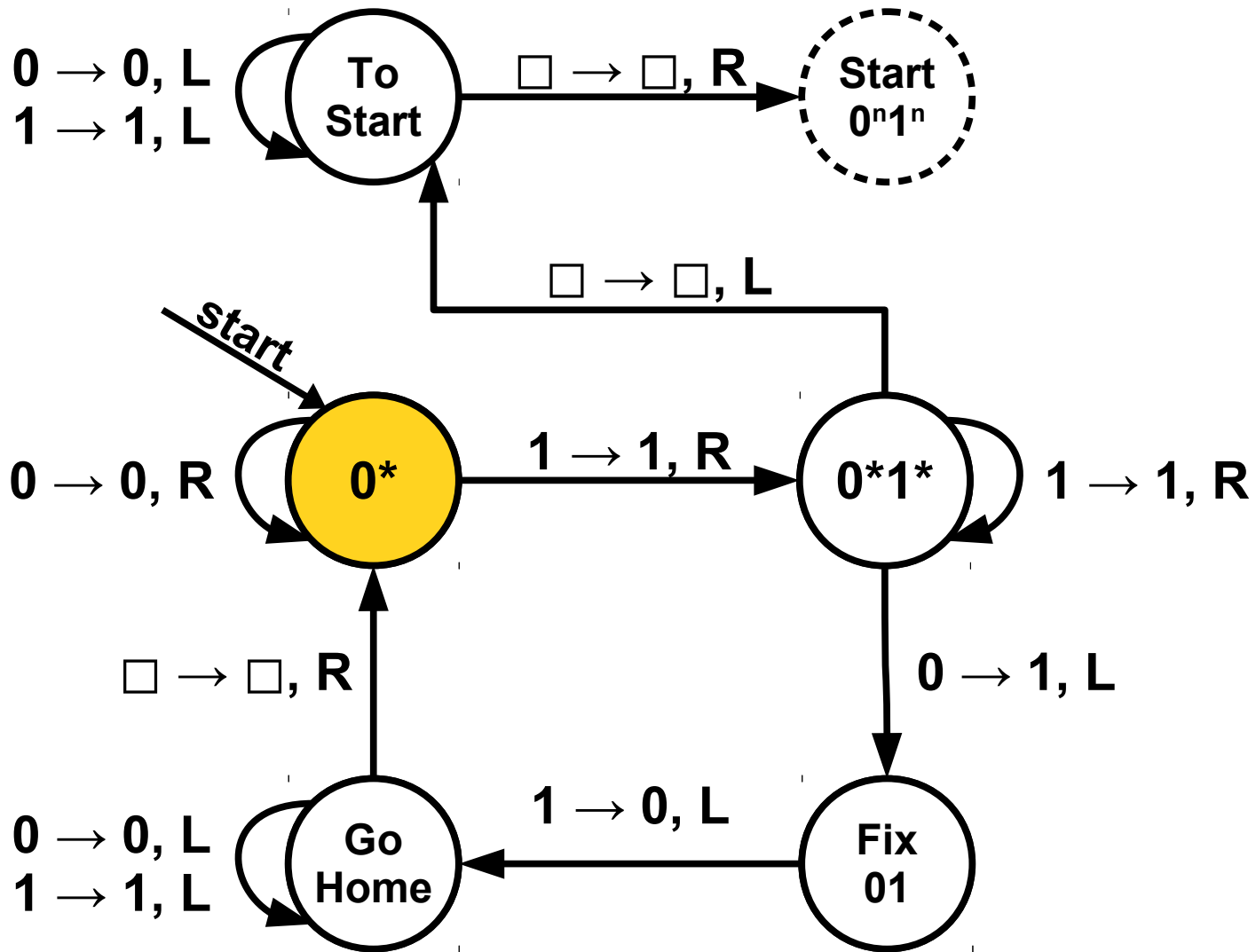
This is just a placeholder. Imagine snapping in the entire TM for  $0^n 1^n$  into this diagram, putting the start state in the dashed area.



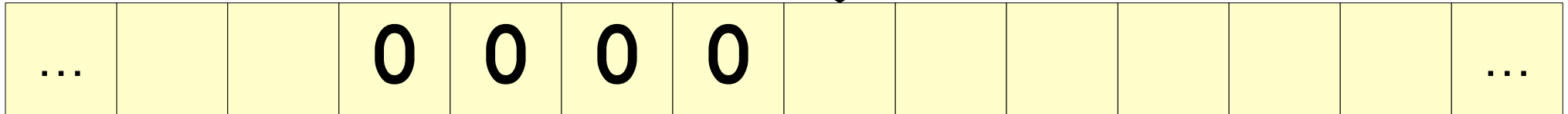
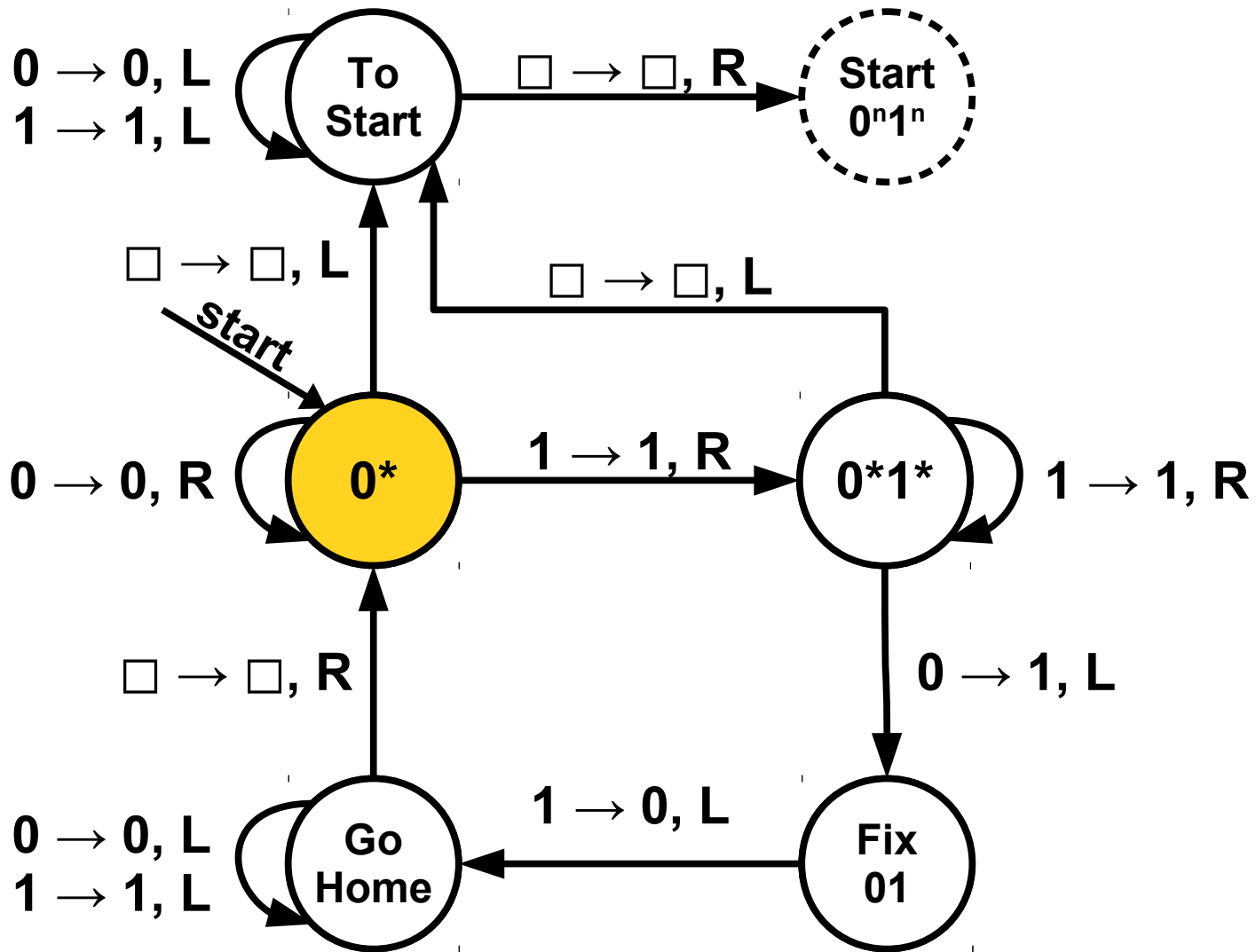






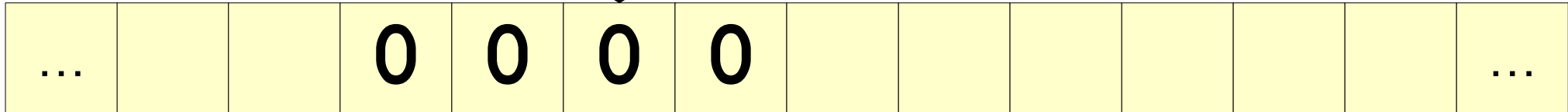
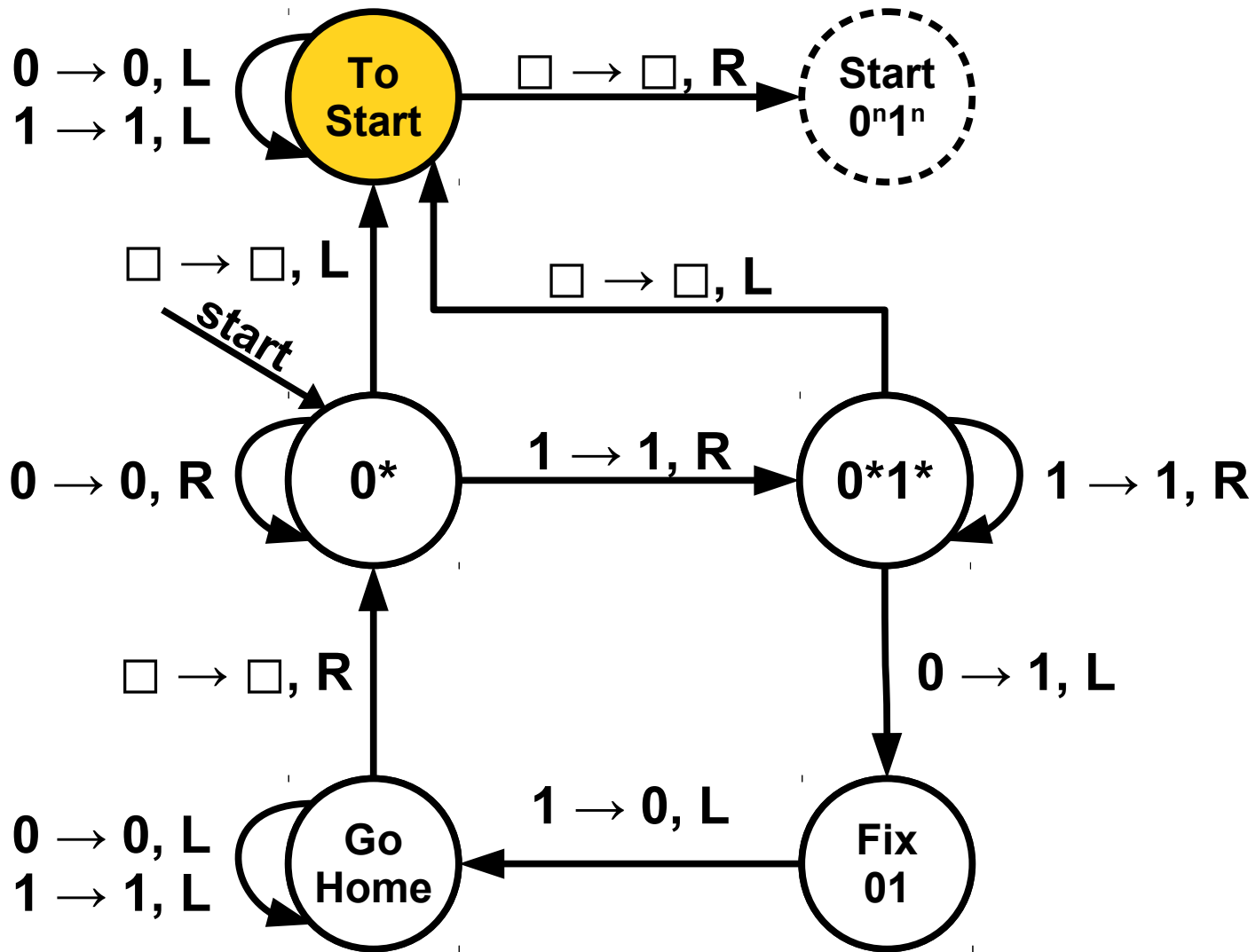




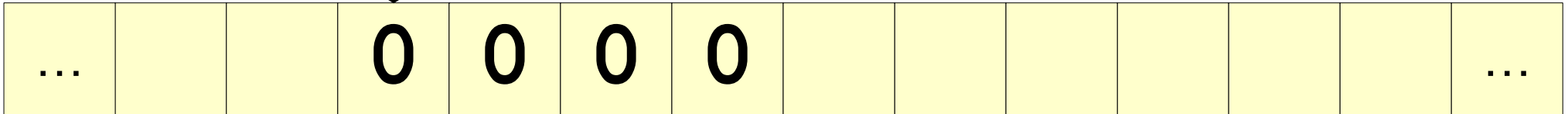
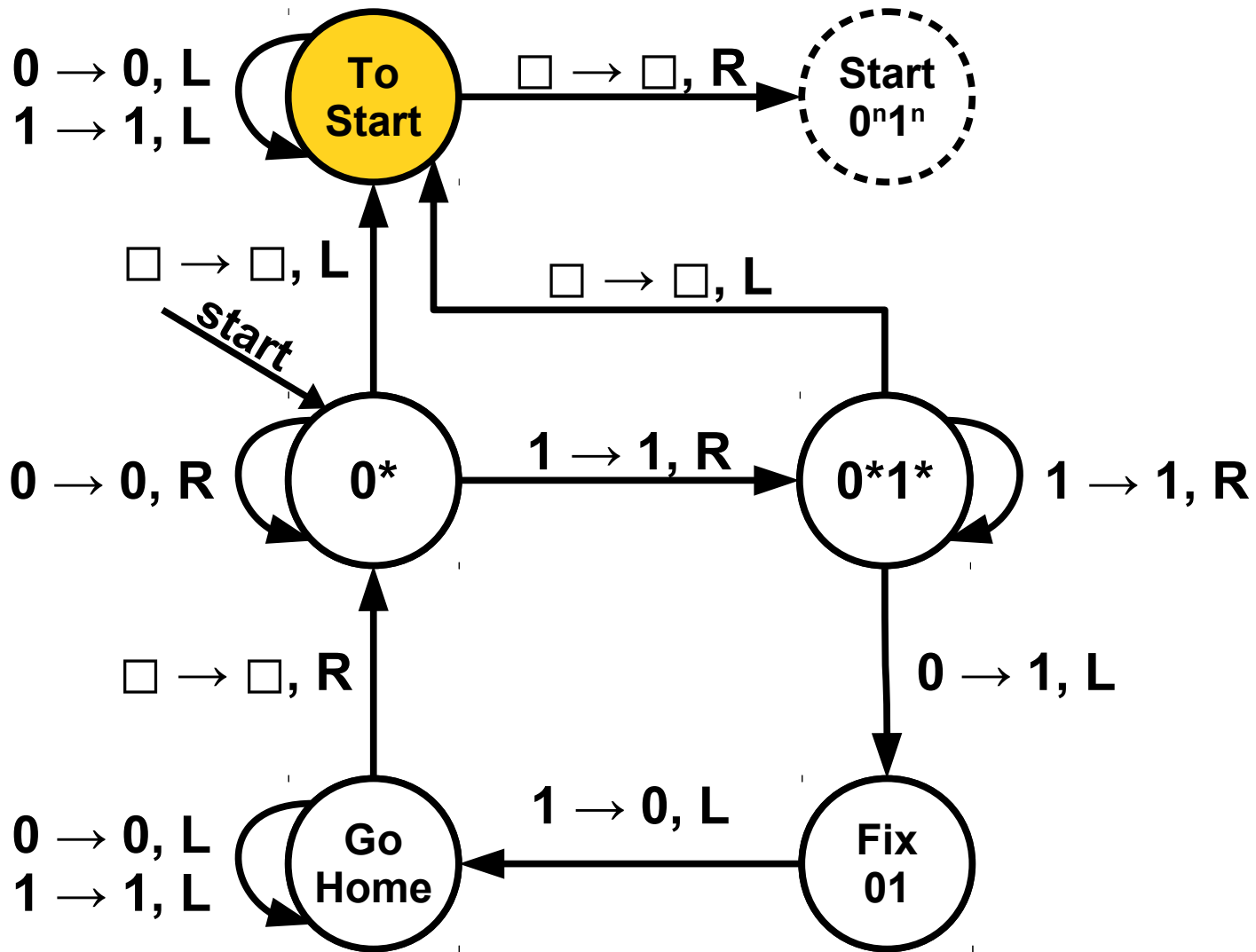




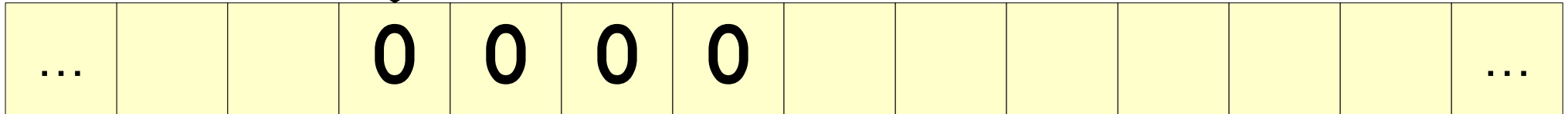
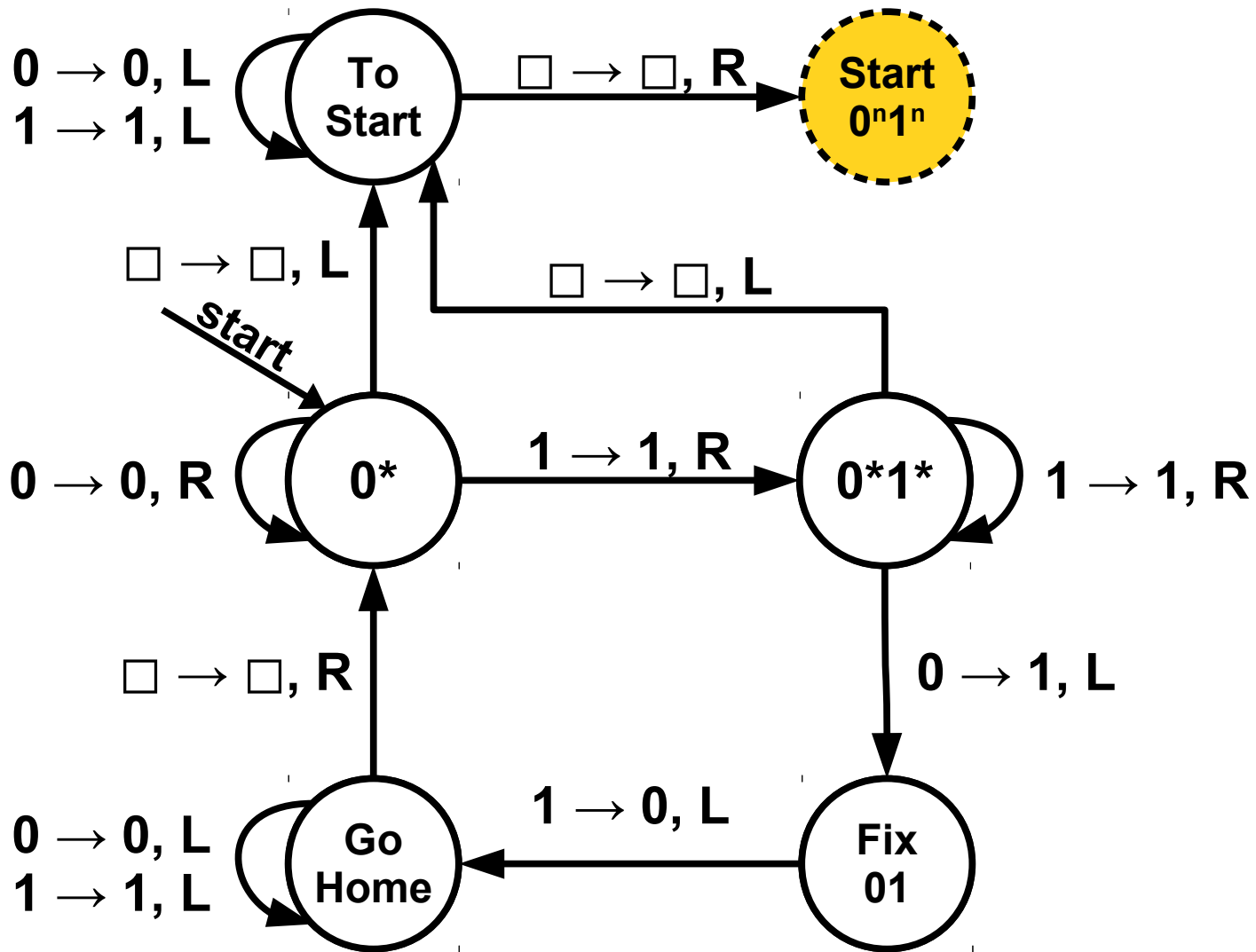


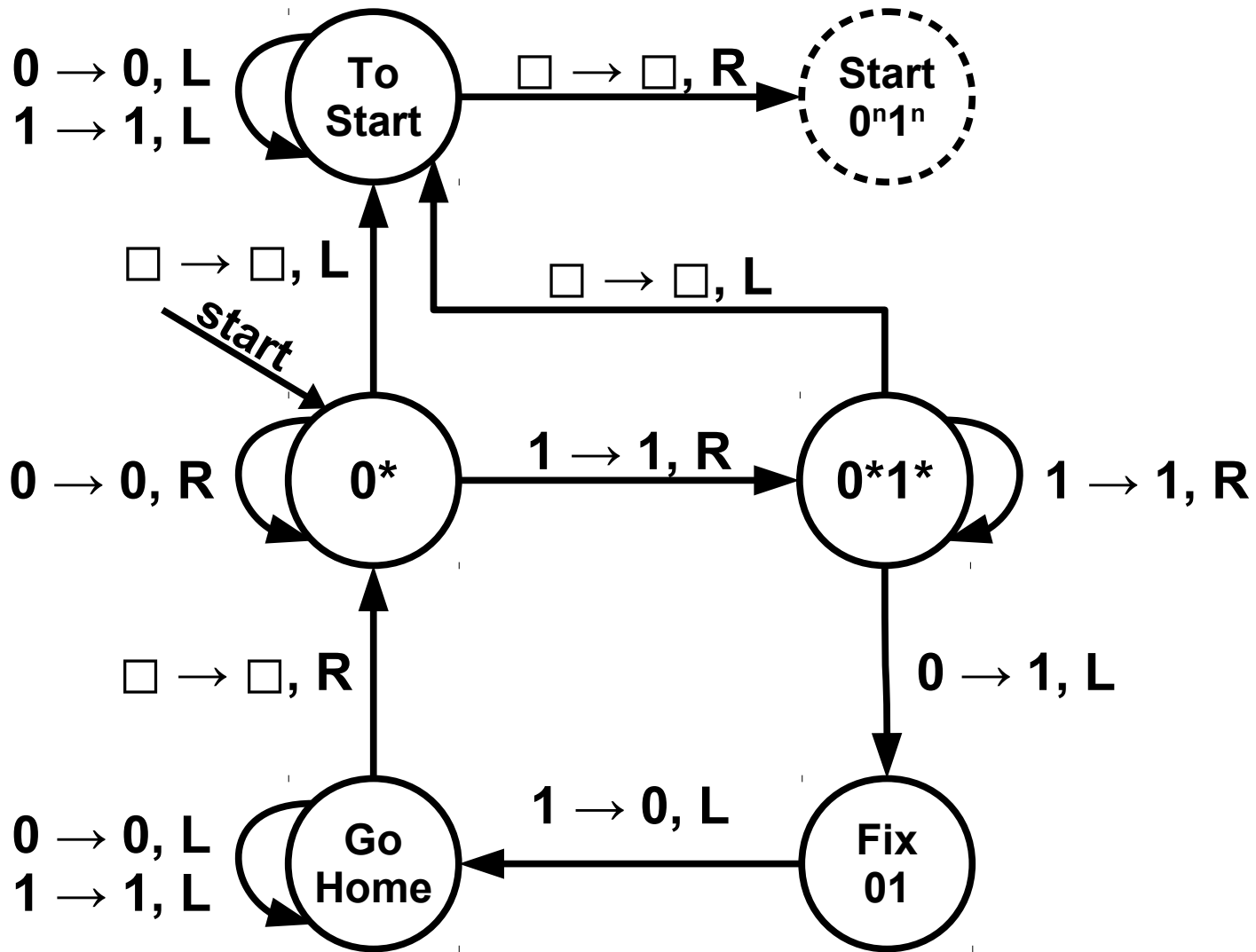


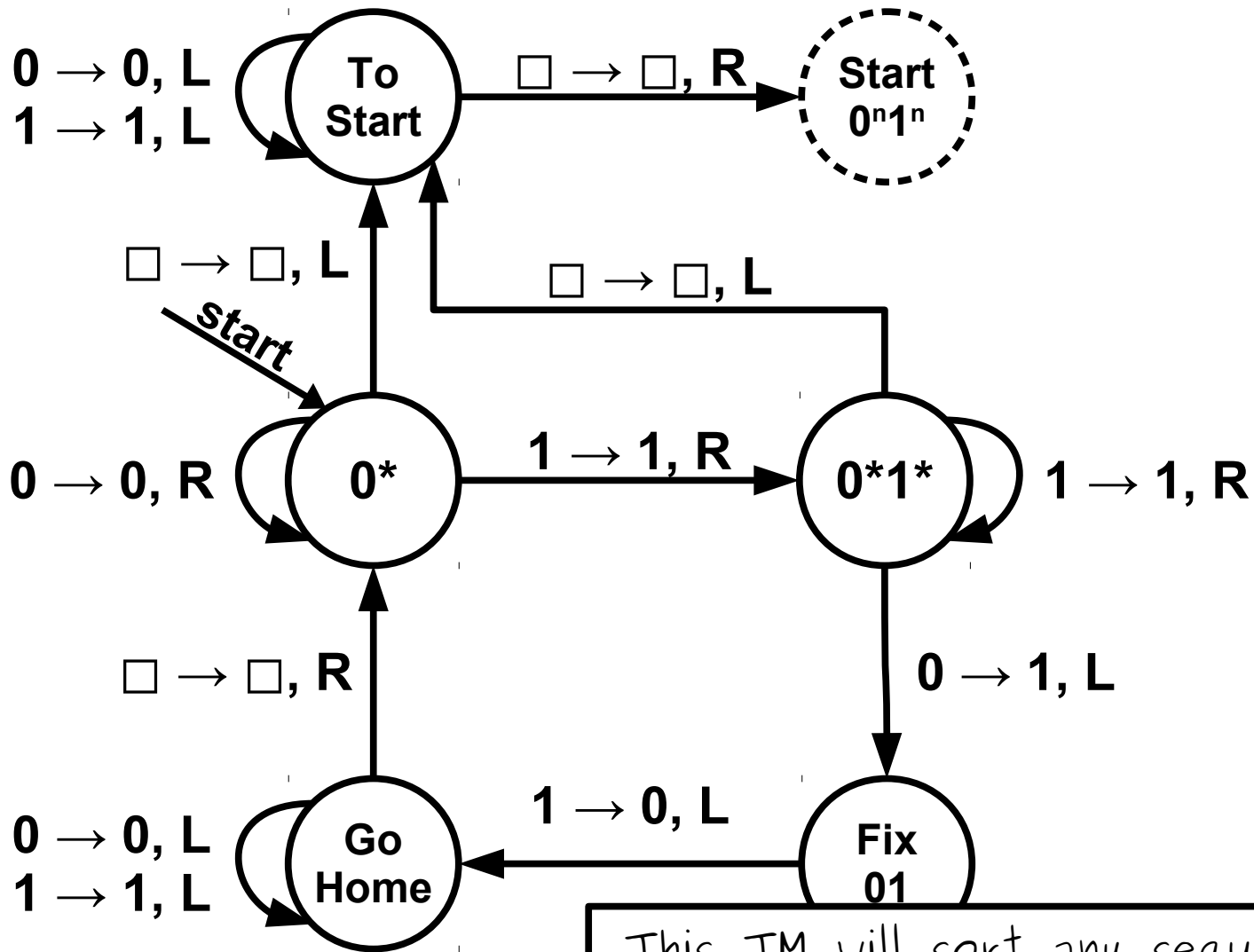












This TM will sort any sequence of 0s and 1s, but it might take a while.

**Fun problem:** design a TM that sorts a string of 0s and 1s, but does so while taking way fewer steps than this machine.



# TM Subroutines

- A ***TM subroutine*** is a Turing machine that, instead of accepting or rejecting an input, does some sort of processing job.
- TM subroutines let us compose larger TMs out of smaller TMs, just as you'd write a larger program using lots of smaller helper functions.
- Here, we saw a TM subroutine that sorts a sequence of 0s and 1s into ascending order.

# TM Subroutines

- Typically, when a subroutine is done running, you have it enter a state marked “done” with a dashed line around it.
- When we're composing multiple subroutines together – which we'll do in a bit – the idea is that we'll snap in some real state for the “done” state.

What other subroutines can we make?