

# Regular Expressions

Recap from Last Time

# Regular Languages

- A language  $L$  is a ***regular language*** if there is a DFA  $D$  such that  $\mathcal{L}(D) = L$ .
- ***Theorem:*** The following are equivalent:
  - $L$  is a regular language.
  - There is a DFA for  $L$ .
  - There is an NFA for  $L$ .

# Language Concatenation

- If  $w \in \Sigma^*$  and  $x \in \Sigma^*$ , then  $wx$  is the **concatenation** of  $w$  and  $x$ .
- If  $L_1$  and  $L_2$  are languages over  $\Sigma$ , the **concatenation** of  $L_1$  and  $L_2$  is the language  $L_1L_2$  defined as

$$L_1L_2 = \{ wx \mid w \in L_1 \text{ and } x \in L_2 \}$$

- Example: if  $L_1 = \{ a, ba, bb \}$  and  $L_2 = \{ aa, bb \}$ , then

$$L_1L_2 = \{ aaa, abb, baaa, babb, bbaa, bbbb \}$$

# Lots and Lots of Concatenation

- Consider the language  $L = \{ \mathbf{aa}, \mathbf{b} \}$
- $LL$  is the set of strings formed by concatenating pairs of strings in  $L$ .

$\{ \mathbf{aaaa}, \mathbf{aab}, \mathbf{baa}, \mathbf{bb} \}$

- $LLL$  is the set of strings formed by concatenating triples of strings in  $L$ .

$\{ \mathbf{aaaaaa}, \mathbf{aaaab}, \mathbf{aabaa}, \mathbf{aabb}, \mathbf{baaaa}, \mathbf{baab}, \mathbf{bbaa}, \mathbf{bbb} \}$

- $LLLL$  is the set of strings formed by concatenating quadruples of strings in  $L$ .

$\{ \mathbf{aaaaaaaa}, \mathbf{aaaaaab}, \mathbf{aaaabaa}, \mathbf{aaaabb}, \mathbf{aabaaaa}, \mathbf{aabaab}, \mathbf{aabbaa}, \mathbf{aabbb}, \mathbf{baaaaaa}, \mathbf{baaaab}, \mathbf{baabaa}, \mathbf{baabb}, \mathbf{bbaaaa}, \mathbf{bbaab}, \mathbf{bbbaa}, \mathbf{bbbb} \}$

# Language Exponentiation

- We can define what it means to “exponentiate” a language as follows:
- $L^0 = \{\varepsilon\}$ 
  - The set containing just the empty string.
  - Idea: Any string formed by concatenating zero strings together is the empty string.
- $L^{n+1} = LL^n$ 
  - Idea: Concatenating  $(n+1)$  strings together works by concatenating  $n$  strings, then concatenating one more.
- **Question:** Why define  $L^0 = \{\varepsilon\}$ ?
- **Question:** What is  $\emptyset^0$ ?

# The Kleene Closure

- An important operation on languages is the ***Kleene Closure***, which is defined as

$$L^* = \{ w \in \Sigma^* \mid \exists n \in \mathbb{N}. w \in L^n \}$$

- Mathematically:

$$w \in L^* \quad \text{iff} \quad \exists n \in \mathbb{N}. w \in L^n$$

- Intuitively, all possible ways of concatenating zero or more strings in  $L$  together, possibly with repetition.
- ***Question:*** What is  $\emptyset^0$ ?

# The Kleene Closure

If  $L = \{ \mathbf{a}, \mathbf{bb} \}$ , then  $L^* = \{$

$\epsilon,$

$\mathbf{a}, \mathbf{bb},$

$\mathbf{aa}, \mathbf{abb}, \mathbf{bba}, \mathbf{bbbb},$

$\mathbf{aaa}, \mathbf{aabb}, \mathbf{abba}, \mathbf{abbbb}, \mathbf{bbaa}, \mathbf{bbabb}, \mathbf{bbbba}, \mathbf{bbbbbb},$

$\dots$

$\}$

Think of  $L^*$  as the set of strings you can make if you have a collection of stamps – one for each string in  $L$  – and you form every possible string that can be made from those stamps.

# Closure Properties

- ***Theorem:*** If  $L_1$  and  $L_2$  are regular languages over an alphabet  $\Sigma$ , then so are the following languages:
  - $\bar{L}_1$
  - $L_1 \cup L_2$
  - $L_1 \cap L_2$
  - $L_1L_2$
  - $L_1^*$
- These properties are called ***closure properties of the regular languages.***

New Stuff!

# Another View of Regular Languages

# Rethinking Regular Languages

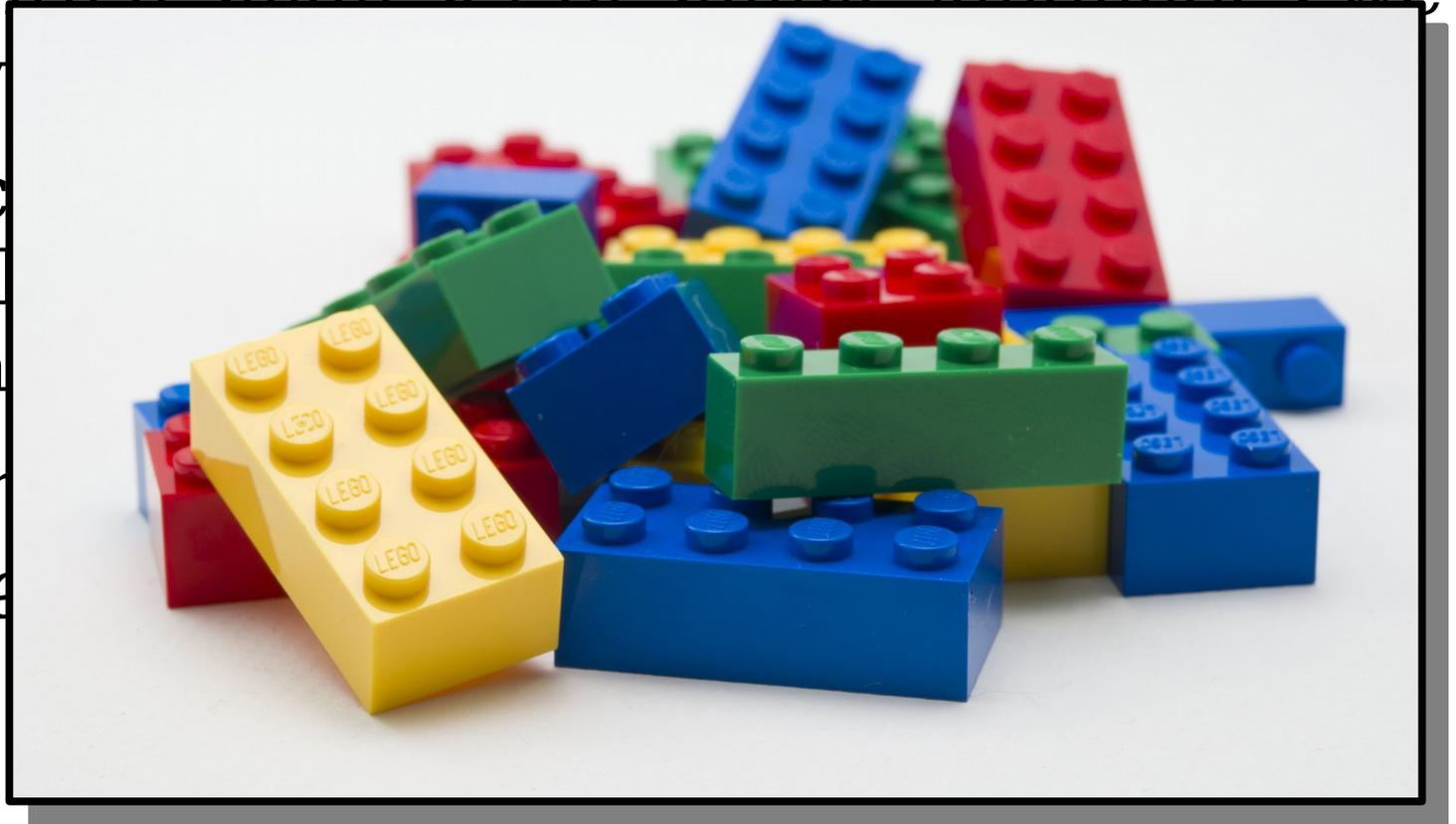
- We currently have several tools for showing a language  $L$  is regular:
  - Construct a DFA for  $L$ .
  - Construct an NFA for  $L$ .
  - Combine several simpler regular languages together via closure properties to form  $L$ .
- We have not spoken much of this last idea.

# Constructing Regular Languages

- **Idea:** Build up all regular languages as follows:
  - Start with a small set of simple languages we already know to be regular.
  - Using closure properties, combine these simple languages together to form more elaborate languages.
- *A bottom-up approach to the regular languages.*

# Constructing Regular Languages

- **Idea:** Build up all regular languages as follows:
  - Start with a small set of simple languages we already know
  - Using closure operations (union, concatenation, Kleene star) to build more complex languages from simple ones
- A bottom-up construction of regular languages



# Regular Expressions

- ***Regular expressions*** are a way of describing a language via a string representation.
- They're used extensively in software systems for string processing and as the basis for tools like grep and flex.
- Conceptually, regular expressions are strings describing how to assemble a larger language out of smaller pieces.

# Atomic Regular Expressions

- The regular expressions begin with three simple building blocks.
- The symbol  $\emptyset$  is a regular expression that represents the empty language  $\emptyset$ .
- For any  $a \in \Sigma$ , the symbol  $a$  is a regular expression for the language  $\{a\}$ .
- The symbol  $\epsilon$  is a regular expression that represents the language  $\{\epsilon\}$ .
  - **Remember:**  $\{\epsilon\} \neq \emptyset!$
  - **Remember:**  $\{\epsilon\} \neq \epsilon!$

# Compound Regular Expressions

- If  $R_1$  and  $R_2$  are regular expressions,  $R_1R_2$  is a regular expression for the *concatenation* of the languages of  $R_1$  and  $R_2$ .
- If  $R_1$  and  $R_2$  are regular expressions,  $R_1 \cup R_2$  is a regular expression for the *union* of the languages of  $R_1$  and  $R_2$ .
- If  $R$  is a regular expression,  $R^*$  is a regular expression for the *Kleene closure* of the language of  $R$ .
- If  $R$  is a regular expression,  $(R)$  is a regular expression with the same meaning as  $R$ .

# Operator Precedence

- Regular expression operator precedence:

$(R)$

$R^*$

$R_1R_2$

$R_1 \cup R_2$

- So **ab\*cUd** is parsed as **((a(b\*))c)Ud**

# Regular Expression Examples

- The regular expression **trickUtreat** represents the regular language { **trick**, **treat** }.
- The regular expression **boo\*** represents the regular language { **boo**, **booo**, **boooo**, ... }.
- The regular expression **candy!(candy!)\*** represents the regular language { **candy!**, **candy!candy!**, **candy!candy!candy!**, ... }.

# Regular Expressions, Formally

- The **language of a regular expression** is the language described by that regular expression.
- Formally:
  - $\mathcal{L}(\epsilon) = \{\epsilon\}$
  - $\mathcal{L}(\emptyset) = \emptyset$
  - $\mathcal{L}(a) = \{a\}$
  - $\mathcal{L}(R_1R_2) = \mathcal{L}(R_1) \mathcal{L}(R_2)$
  - $\mathcal{L}(R_1 \cup R_2) = \mathcal{L}(R_1) \cup \mathcal{L}(R_2)$
  - $\mathcal{L}(R^*) = \mathcal{L}(R)^*$
  - $\mathcal{L}((R)) = \mathcal{L}(R)$

Worthwhile activity: Apply this recursive definition to

**a(bUc)((d))**

and see what you get.

# Designing Regular Expressions

- Let  $\Sigma = \{a, b\}$ .
- Let  $L = \{ w \in \Sigma^* \mid w \text{ contains } aa \text{ as a substring} \}$ .

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**bbabbbaabab**

**aaaa**

**bbbbbabbbbaabbbb**

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**aaa**

bbbbbabbbb**aa**bbbbbb

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$\Sigma^*aa\Sigma^*$

bbabbb**aa**bab

**aaa**

bbbbbabbbb**aa**bbbbbb

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The length of  
a string  $w$  is  
denoted  $|w|$

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$\Sigma\Sigma\Sigma\Sigma$

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$\Sigma\Sigma\Sigma\Sigma$

**aaaa**

**baba**

**bbbb**

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$\Sigma\Sigma\Sigma\Sigma$

aaaa

baba

bbbb

baaa

# Designing Regular Expressions

- Let  $\Sigma = \{a, b\}$ .
- Let  $L = \{ w \in \Sigma^* \mid |w| = 4 \}$ .

$\Sigma^4$

aaaa  
baba  
bbbb  
baaa

# Designing Regular Expressions

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$\Sigma^4$

aaaa  
baba  
bbbb  
baaa

# Designing Regular Expressions

- Let  $\Sigma = \{a, b\}$ .
- Let  $L = \{ w \in \Sigma^* \mid w \text{ contains at most one } a \}$ .

Here are some candidate regular expressions for the language  $L$ . Which of these are correct?

$\Sigma^*a\Sigma^*$

$b^*ab^* \cup b^*$

$b^*(a \cup \epsilon)b^*$

$b^*a^*b^* \cup b^*$

$b^*(a^* \cup \epsilon)b^*$

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abbb

a

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bbbbabbb

bbbbbb

abbb

a

# Designing Regular Expressions

- Let  $\Sigma = \{a, b\}$ .
- Let  $L = \{ w \in \Sigma^* \mid w \text{ contains at most one } a \}$ .

$b^*a?b^*$

bbbbabbb

bbbbbb

abbb

a

# A More Elaborate Design

- Let  $\Sigma = \{ \mathbf{a}, ., @ \}$ , where  $\mathbf{a}$  represents “some letter.”
- Let's make a regex for email addresses.

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**a**<sup>+</sup> (.aa\*)\* @ aa\*.aa\* (.aa\*)\*

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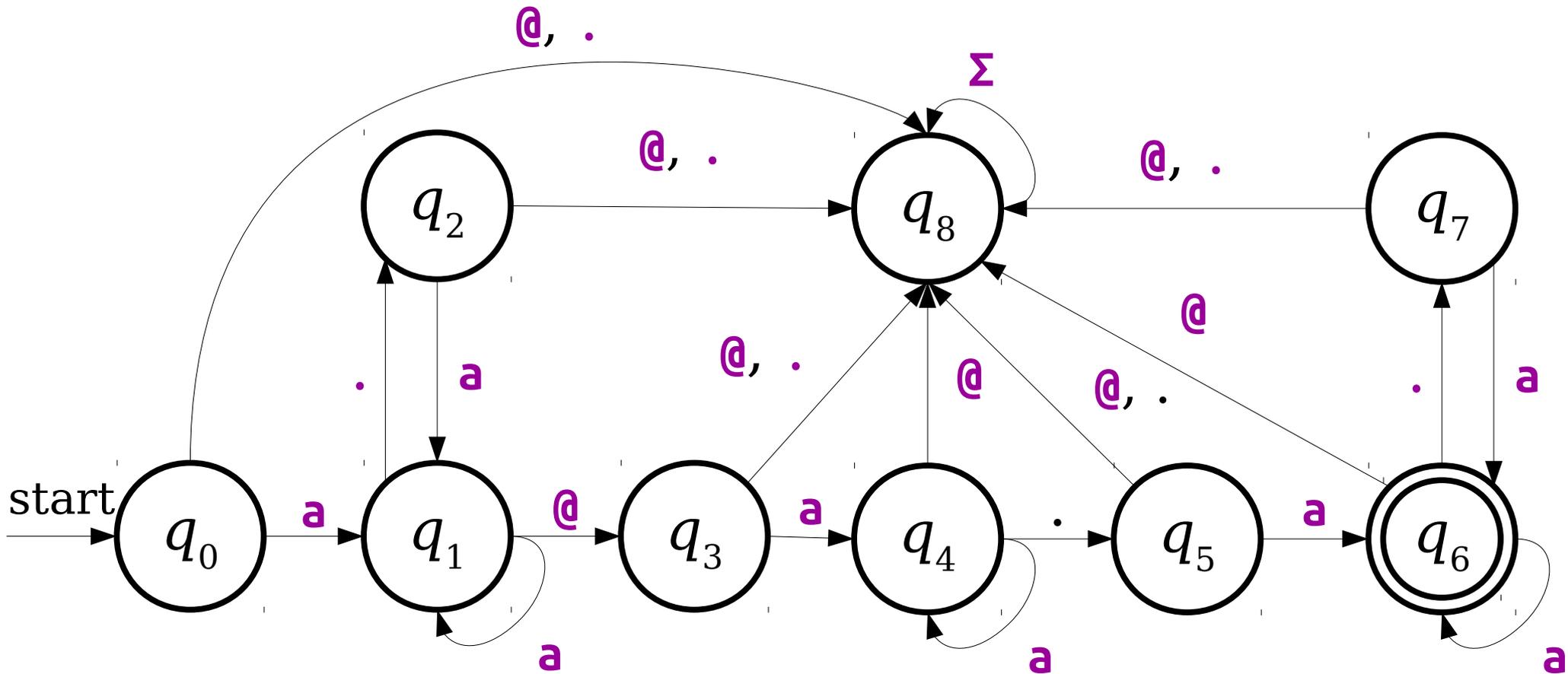
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$\mathbf{a}^+ (\mathbf{.a}^+)^* \mathbf{@} \mathbf{a}^+ (\mathbf{.a}^+)^+$

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# For Comparison

$a^+ (.a^+) * @a^+ (.a^+)^+$



# Shorthand Summary

- $R^n$  is shorthand for  $RR \dots R$  ( $n$  times).
  - Edge case: define  $R^0 = \varepsilon$ .
- $\Sigma$  is shorthand for “any character in  $\Sigma$ .”
- $R?$  is shorthand for  $(R \cup \varepsilon)$ , meaning “zero or one copies of  $R$ .”
- $R^+$  is shorthand for  $RR^*$ , meaning “one or more copies of  $R$ .”

**Time-Out for Announcements!**

Stanford Women  
in Computer Science

# CASUAL DINNER



Tuesday, November 6th from 5-7 PM at Gates  
403

Come have dinner with CS students and faculty.  
Everyone is welcome, especially students  
just starting out in CS!

# Engineers for a Sustainable World

- Engineers for a Sustainable World (ESW) is holding an information session this Thursday, November 8, at 8:00PM in the Crothers Hall Big Lounge.
- There are some *very* cool projects you can get involved with here involving power grids and archaeology.
- Want to attend? RSVP using [this link](#).

# Cardinal Quarter Opportunities Fair

- Want to work on a public service project?
- Want over 400 choices to pick from?
- Stop by the Cardinal Quarter Opportunities Fair on Tuesday, November 13<sup>th</sup>.
  - 6:30PM – 8:00PM, Tresidder Union, Oak Lounge.

# Midterm Exam Logistics

- The next midterm is **Monday, November 12<sup>th</sup>** from **7:00PM - 10:00PM**. Locations are divvied up by last (family) name:
  - A-L: Go to **Bishop Auditorium**.
  - M-Z: Go to **Cubberley Auditorium**.
- The exam focuses on Lecture 06 – 13 (binary relations through induction, inclusive) and PS3 – PS5. Finite automata onward is *not* tested.
  - Topics from earlier in the quarter (proofwriting, first-order logic, set theory, etc.) are also fair game, but that's primarily because the later material builds on this earlier material.
- The exam is closed-book, closed-computer, and limited-note. You can bring a double-sided, 8.5" × 11" sheet of notes with you to the exam, decorated however you'd like.
- Students with OAE accommodations: please contact us **immediately** if you haven't yet done so. We'll ping you about setting up alternate exams.

# Practice Midterm Exam

- We'll be holding a practice midterm exam **Wednesday** from **7PM - 10PM** in **320-105**.
- The practice midterm exam is composed of what we think is a good representative sample of older midterm questions from across the years. It's probably the best indicator of what you should expect to see.
- Course staff will be on hand to answer your questions.
- Can't make it? We'll release the practice exam and solutions online. Set up your own practice exam time with a small group and work through it under realistic conditions!

# Other Practice Materials

- We've posted five practice midterms to the course website, with solutions.
  - We'll post the practice exam from this evening a little bit later, bringing the total to six.
- There's also Extra Practice Problems 2, plus all the CS103A materials.
- Need more practice? Let us know and we'll see what we can do!

# Problem Sets

- Problem Set Five solutions are now out.
  - ***Please read them!*** We've included long, detailed explanations for some of the trickier problems that show off how to get to the solution.
  - We'll aim to have PS5 graded and returned on Wednesday evening.
- Problem Set Six is out and is due this Friday at 2:30PM.
  - ***Be careful about using late days here***, since the exam is on Monday.

Your Questions

“Is it okay to end a friendship that’s become too emotionally draining even if the other person could really use your support?”

I'll take this one in class because it doesn't have a short answer.

Back to CS103!

# The Power of Regular Expressions

***Theorem:*** If  $R$  is a regular expression, then  $\mathcal{L}(R)$  is regular.

***Proof idea:*** Use induction!

- The atomic regular expressions all represent regular languages.
- The combination steps represent closure properties.
- So anything you can make from them must be regular!

# Thompson's Algorithm

- In practice, many regex matchers use an algorithm called ***Thompson's algorithm*** to convert regular expressions into NFAs (and, from there, to DFAs).
  - Read Sipser if you're curious!
- ***Fun fact:*** the “Thompson” here is Ken Thompson, one of the co-inventors of Unix!

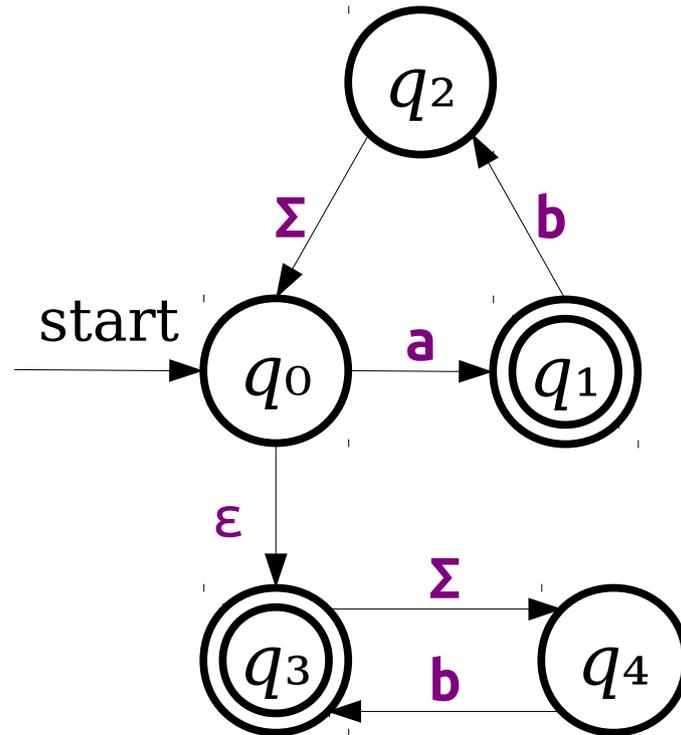
# The Power of Regular Expressions

***Theorem:*** If  $L$  is a regular language, then there is a regular expression for  $L$ .

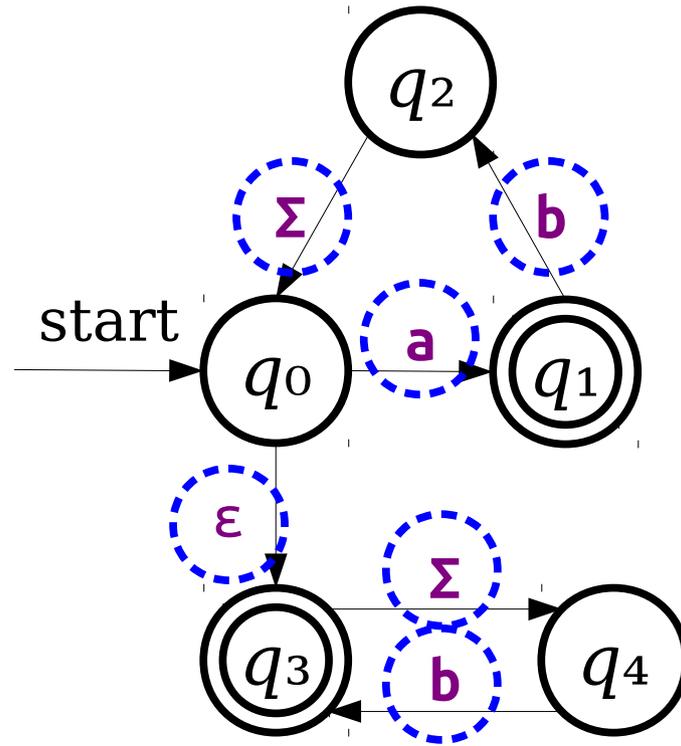
***This is not obvious!***

***Proof idea:*** Show how to convert an arbitrary NFA into a regular expression.

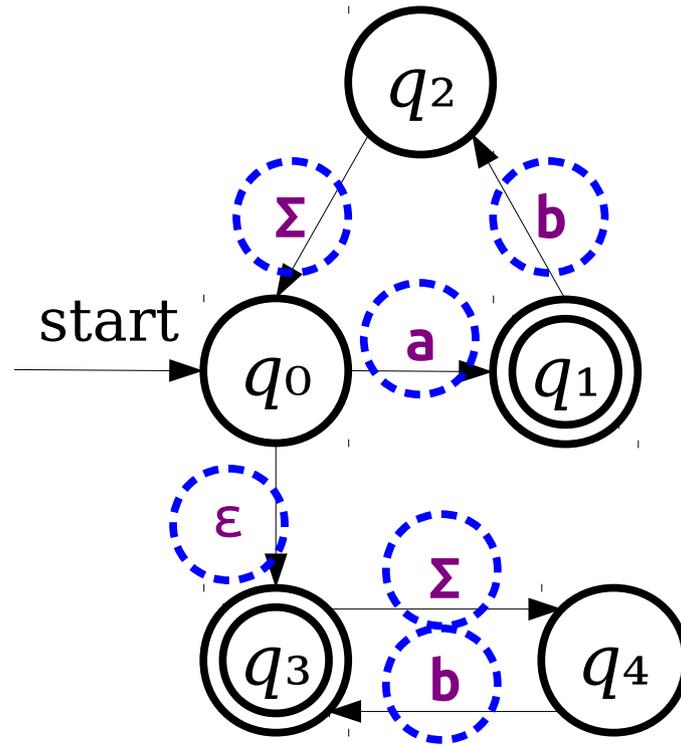
# Generalizing NFAs



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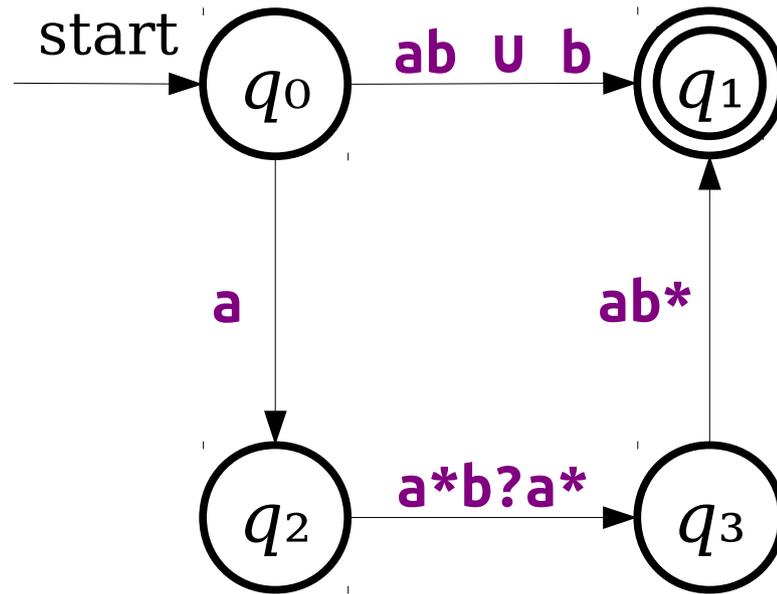


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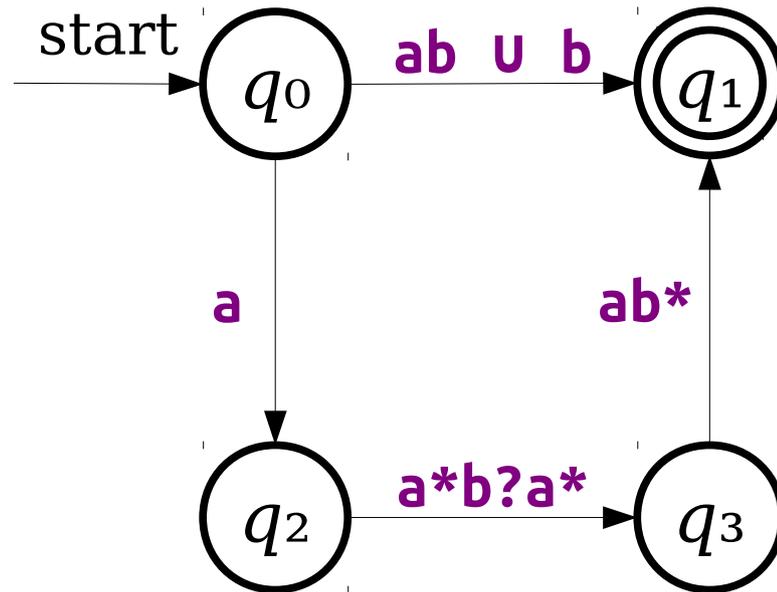


These are all regular expressions!

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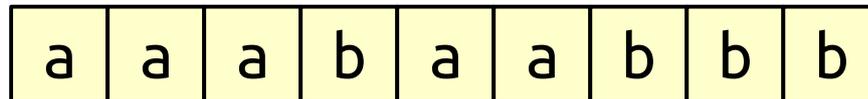
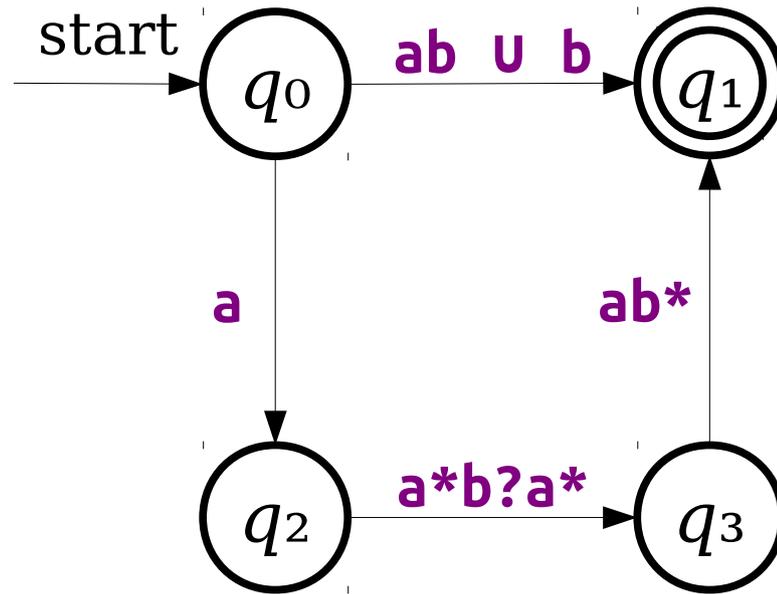


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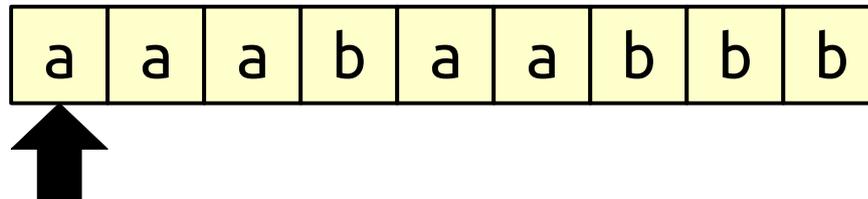
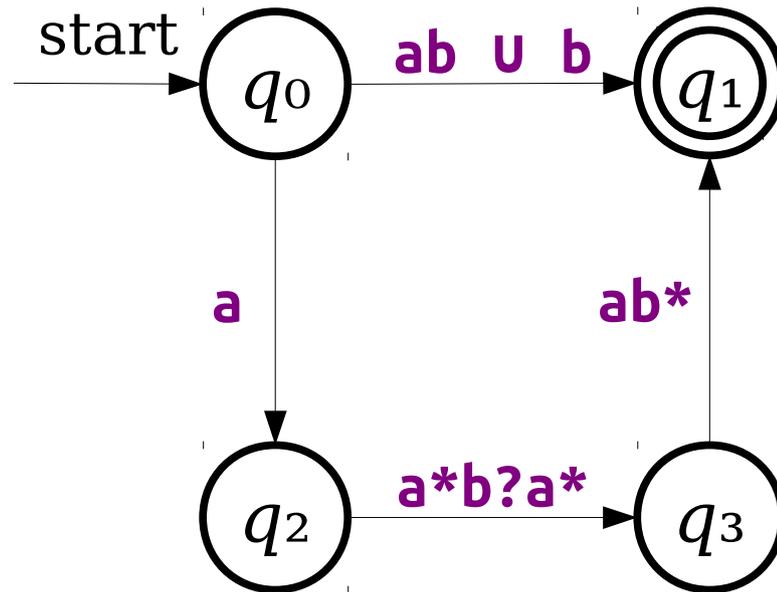


Note: Actual NFAs aren't allowed to have transitions like these. This is just a thought experiment.

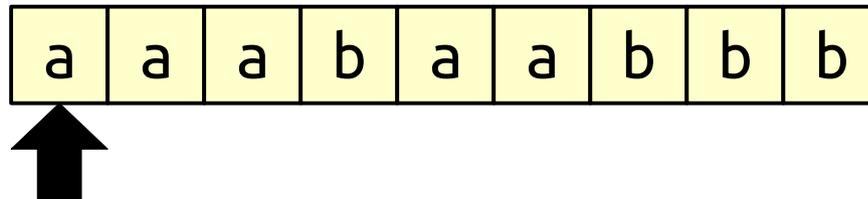
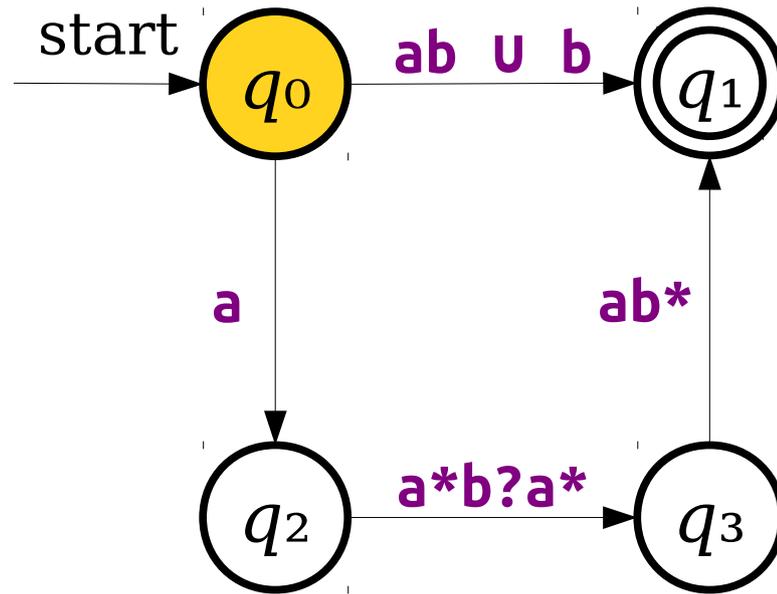
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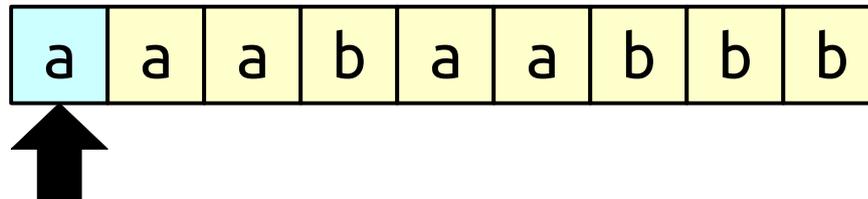
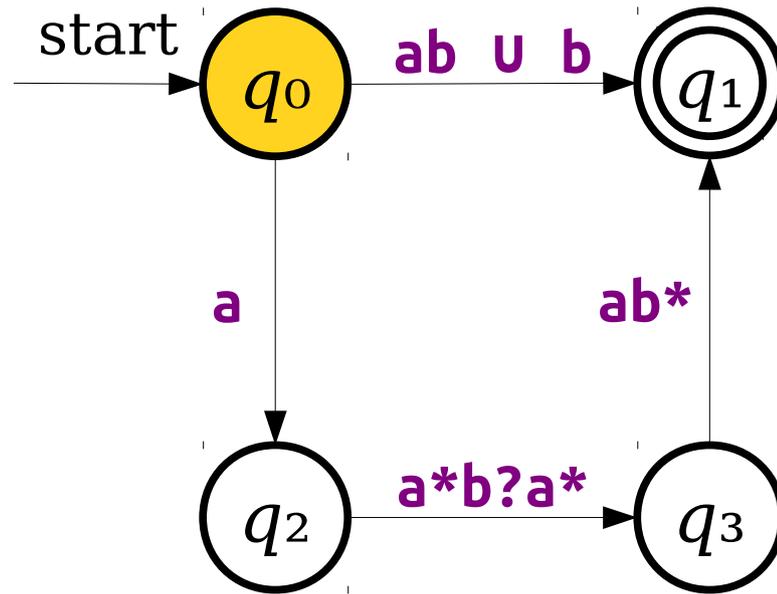
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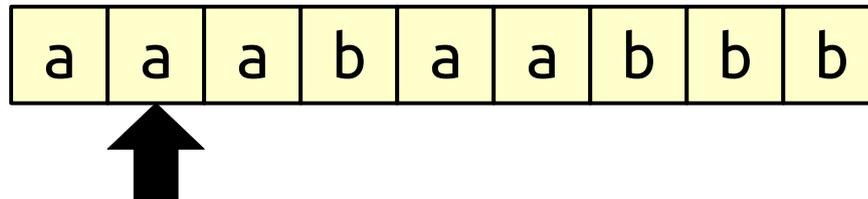
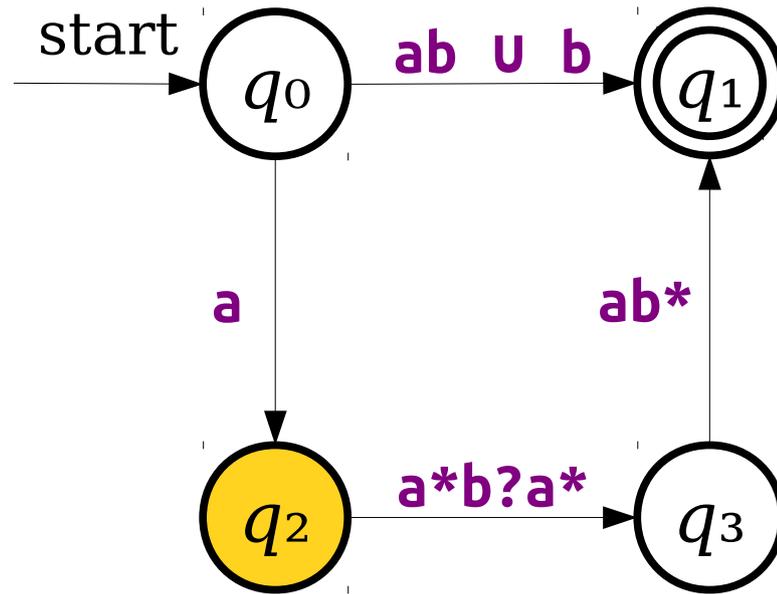
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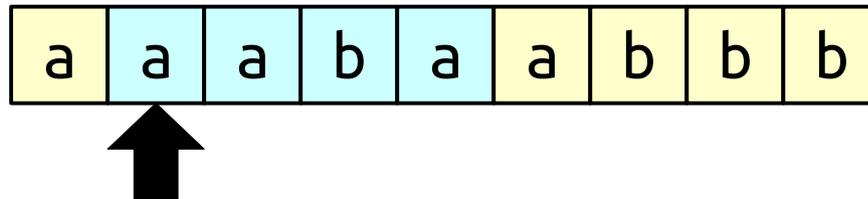
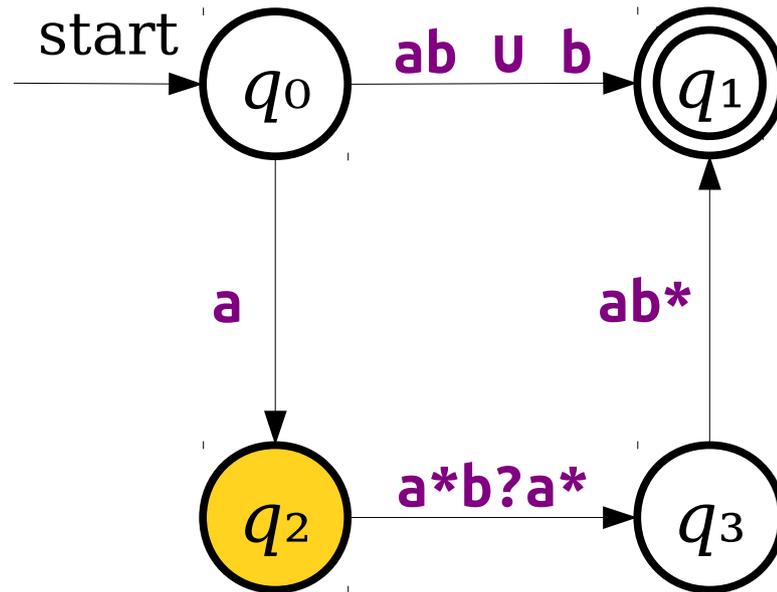
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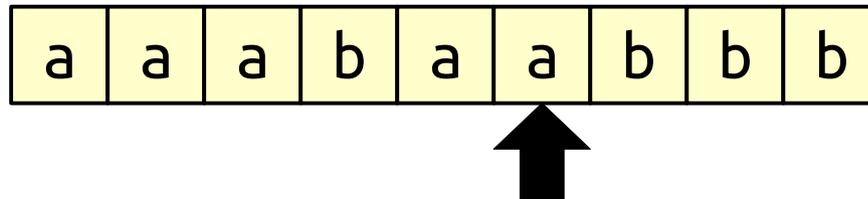
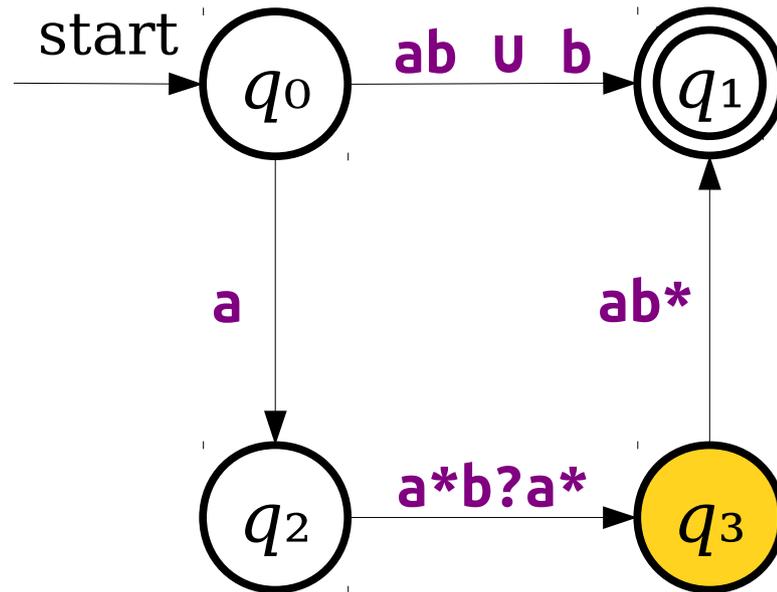
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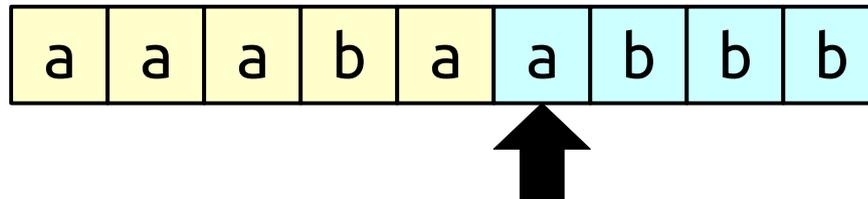
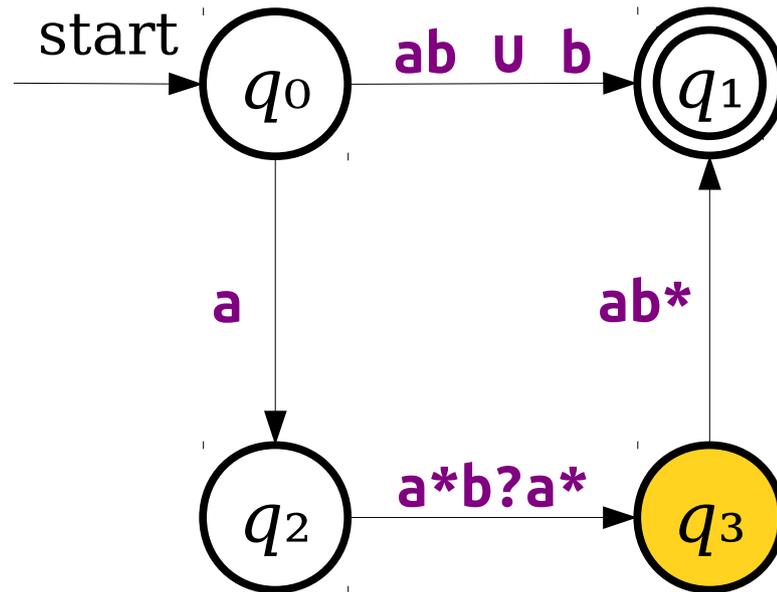
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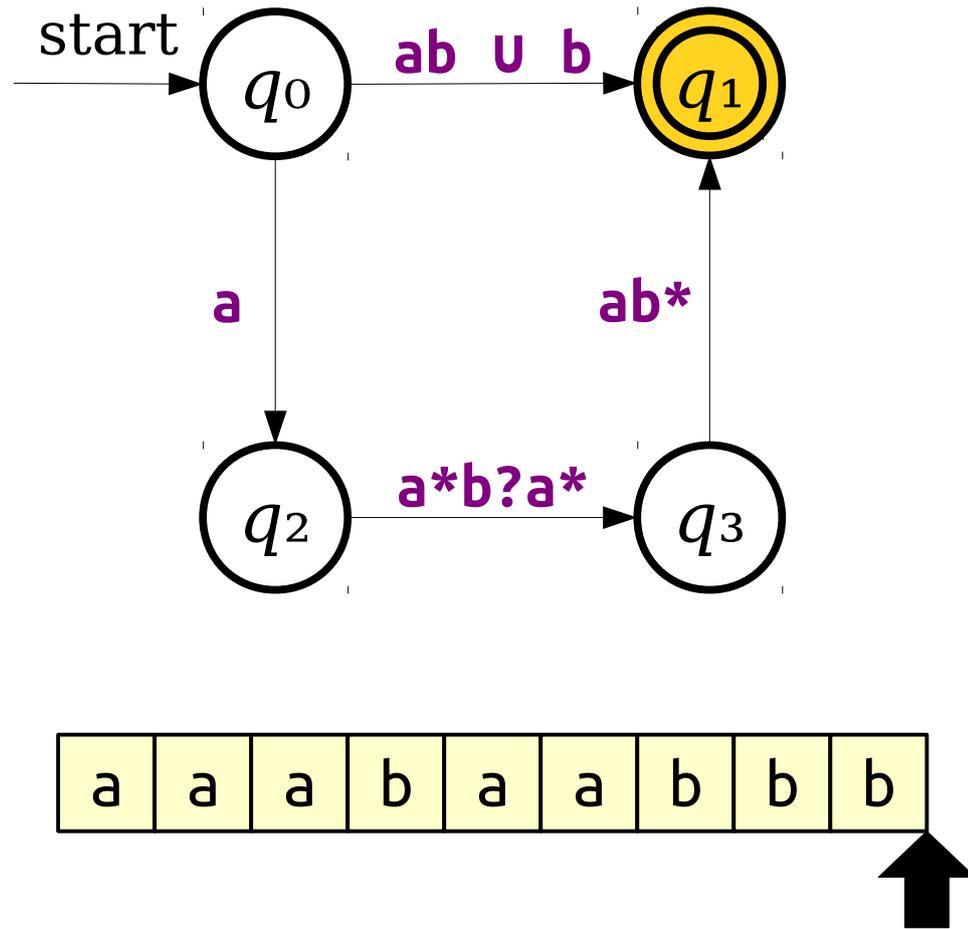
# Generalizing NFAs



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***Key Idea 1:*** Imagine that we can label transitions in an NFA with arbitrary regular expressions.

# Generalizing NFAs

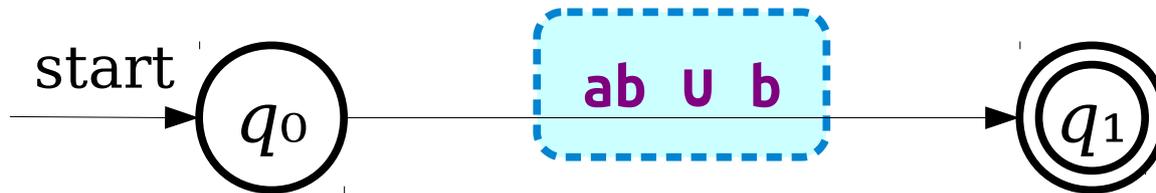


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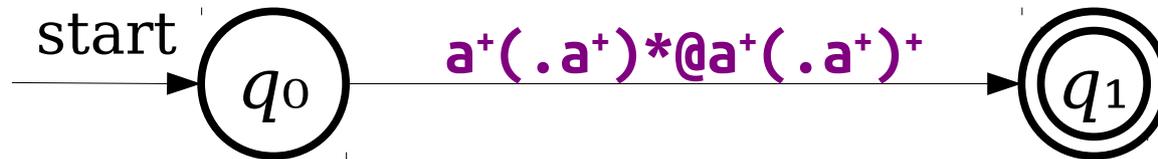
Is there a simple regular expression for the language of this generalized NFA?

# Generalizing NFAs

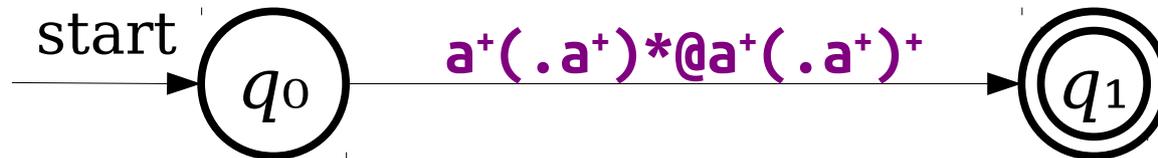


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# Generalizing NFAs

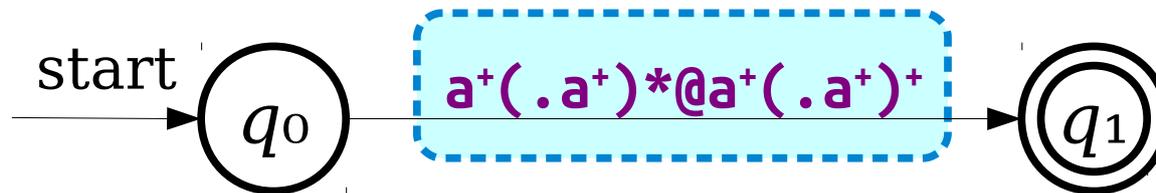


# Generalizing NFAs



Is there a simple regular expression for the language of this generalized NFA?

# Generalizing NFAs



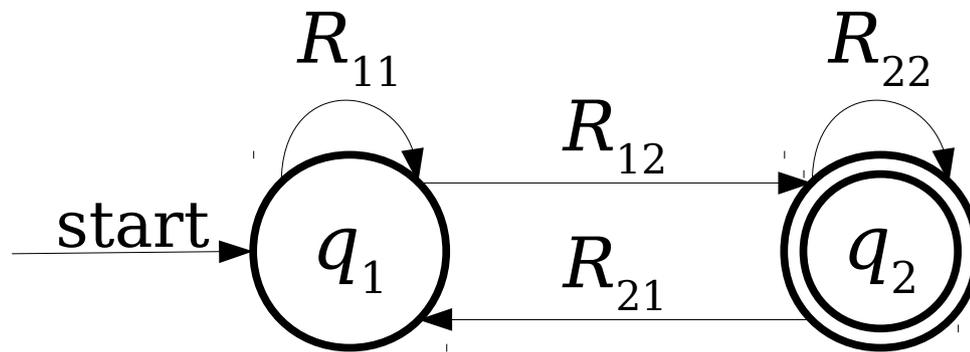
Is there a simple regular expression for the language of this generalized NFA?

**Key Idea 2:** If we can convert an NFA into a generalized NFA that looks like this...

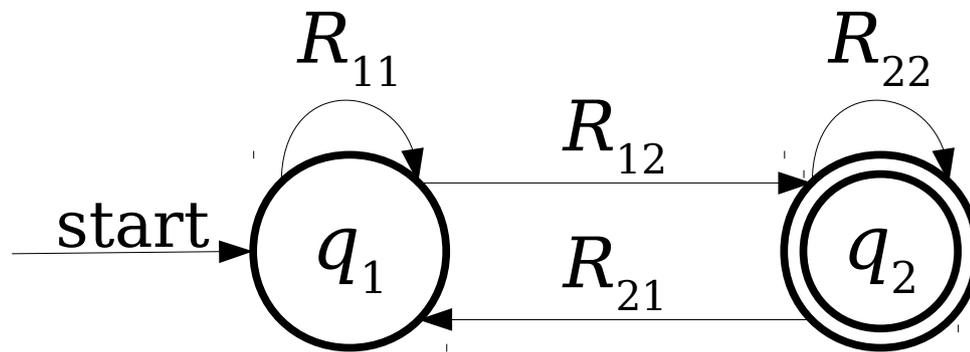


...then we can easily read off a regular expression for the original NFA.

# From NFAs to Regular Expressions

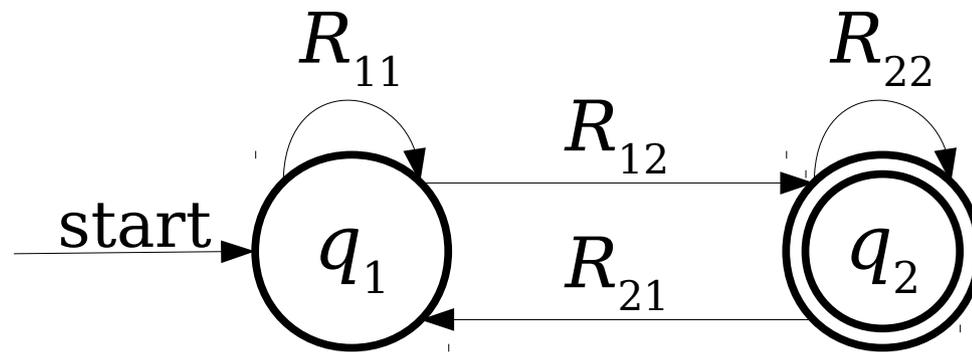


# From NFAs to Regular Expressions



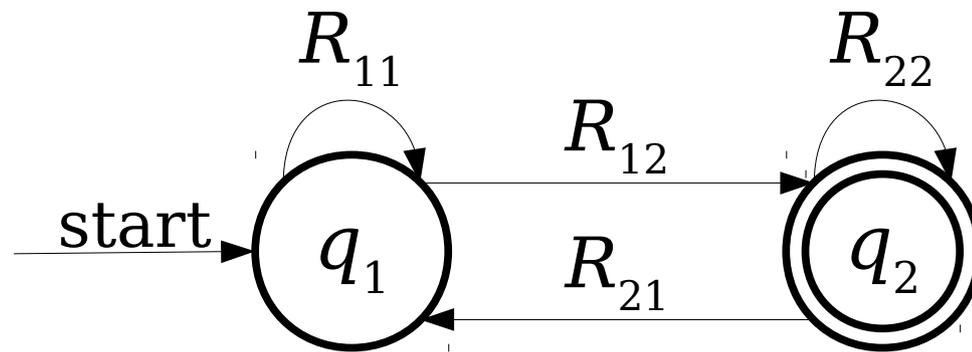
Here,  $R_{11}$ ,  $R_{12}$ ,  $R_{21}$ , and  $R_{22}$  are arbitrary regular expressions.

# From NFAs to Regular Expressions



Question: Can we get a clean regular expression from this NFA?

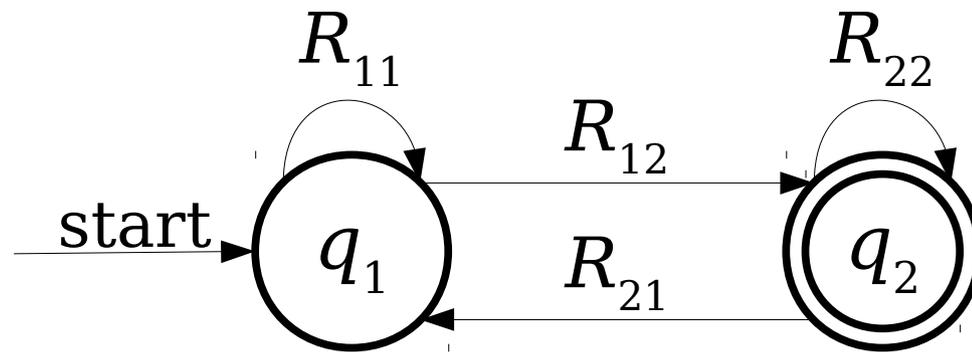
# From NFAs to Regular Expressions



Key Idea 3: Somehow transform this NFA so that it looks like this:

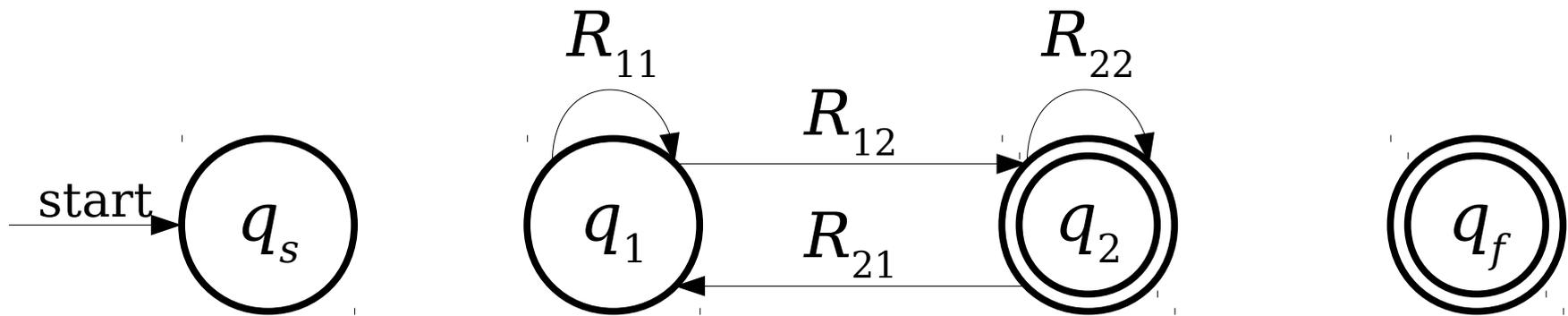


# From NFAs to Regular Expressions

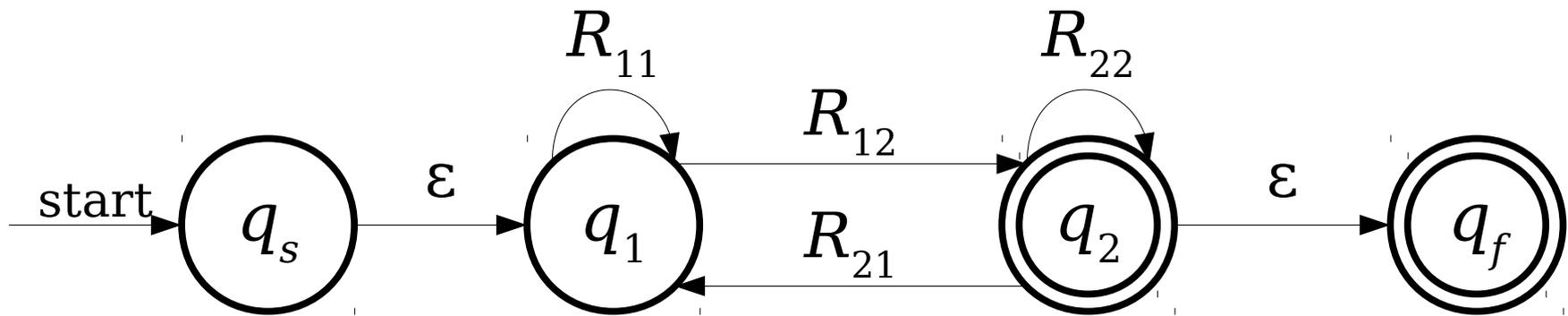


The first step is going to be a bit weird...

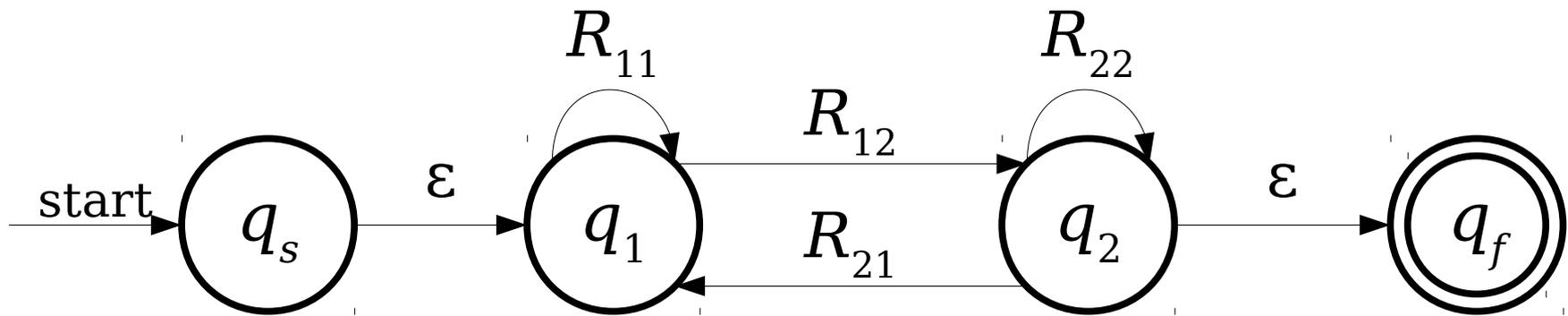
# From NFAs to Regular Expressions



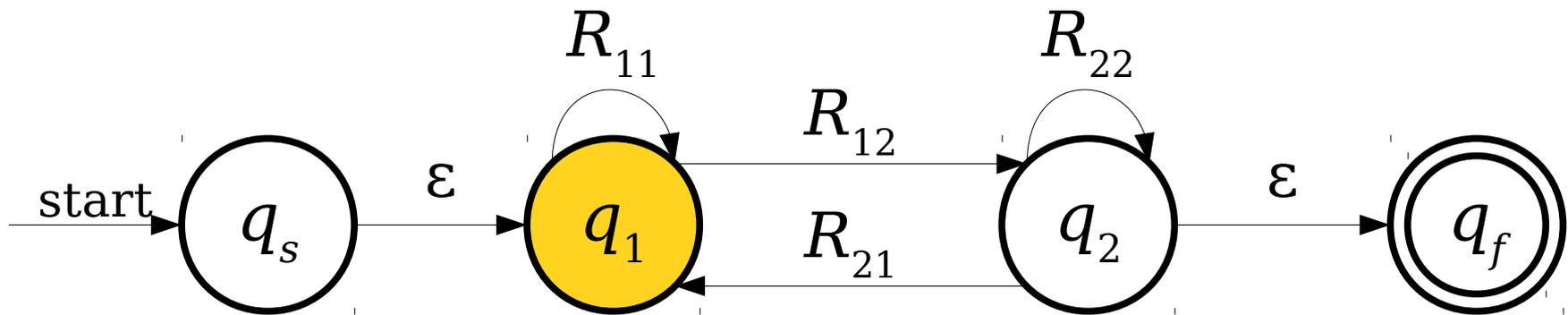
# From NFAs to Regular Expressions



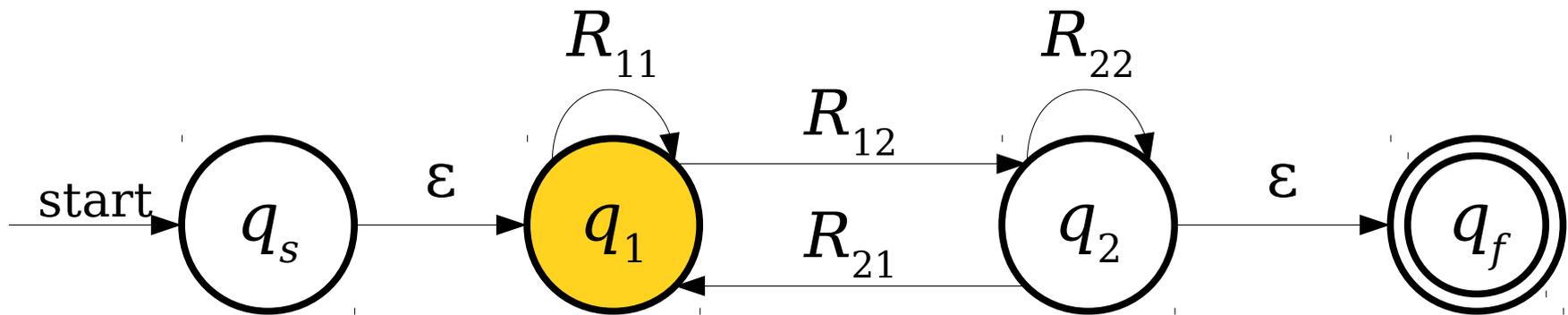
# From NFAs to Regular Expressions



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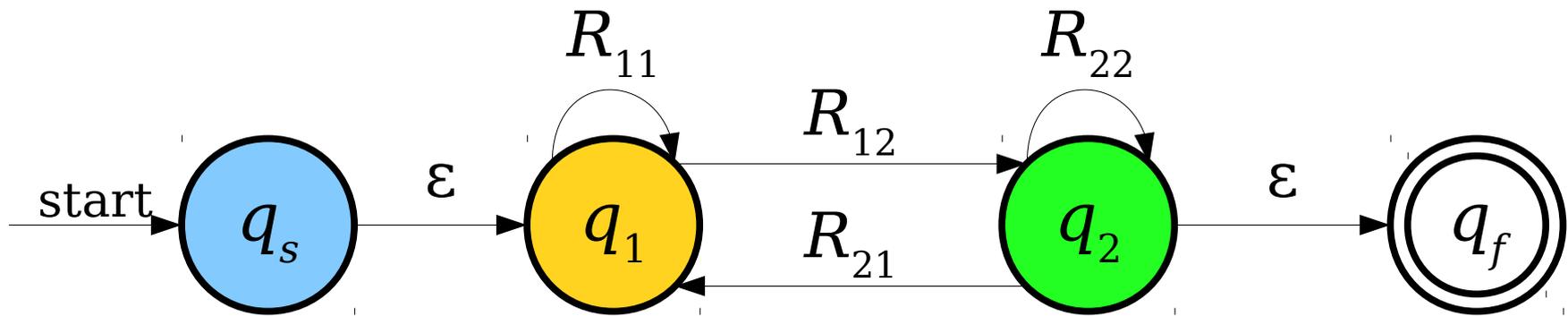


# From NFAs to Regular Expressions

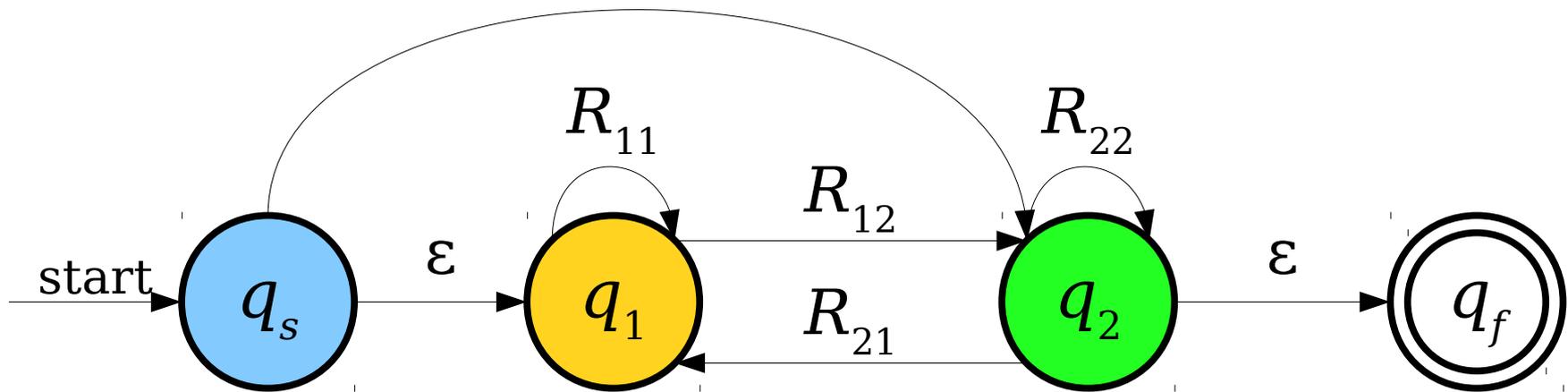


Could we eliminate  
this state from  
the NFA?

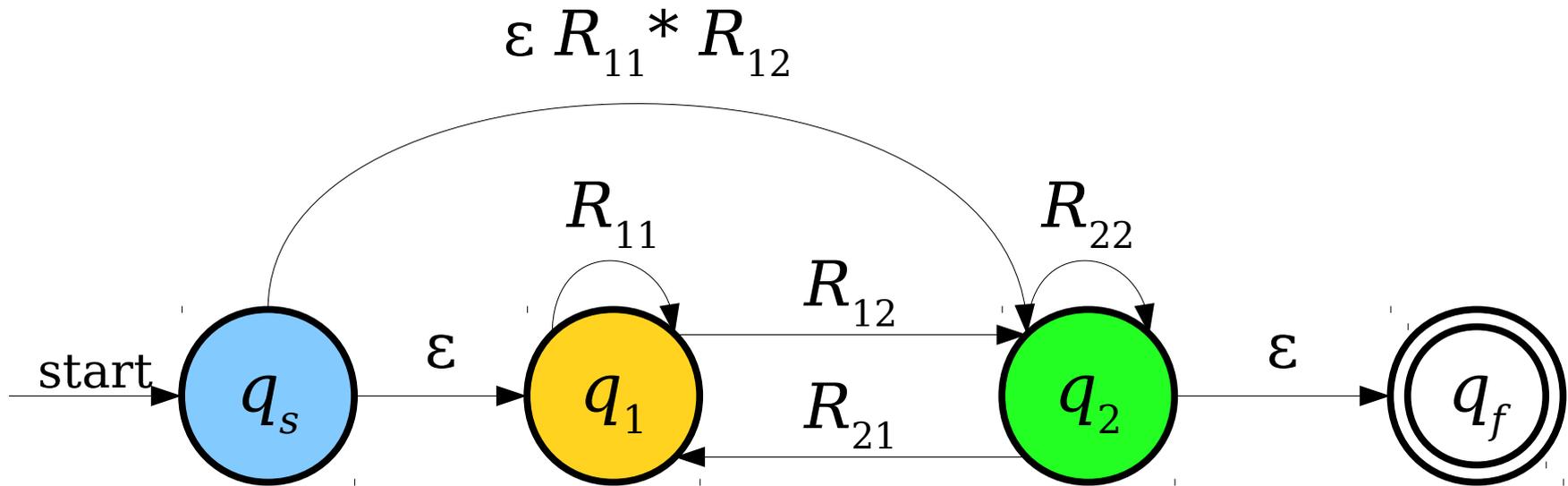
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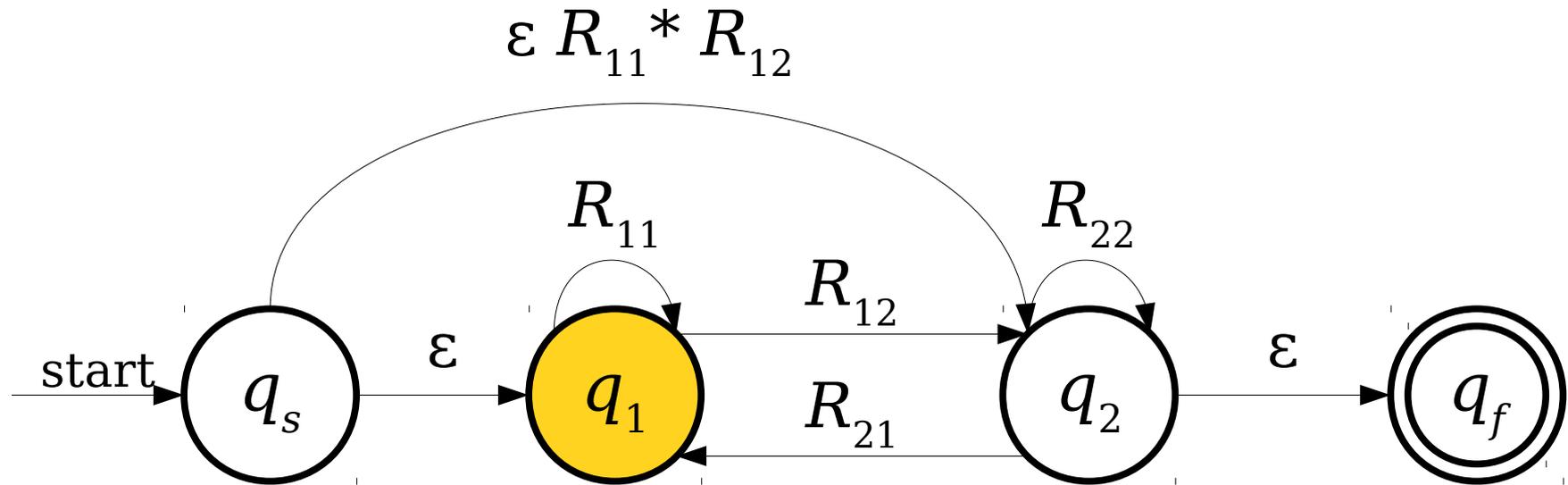


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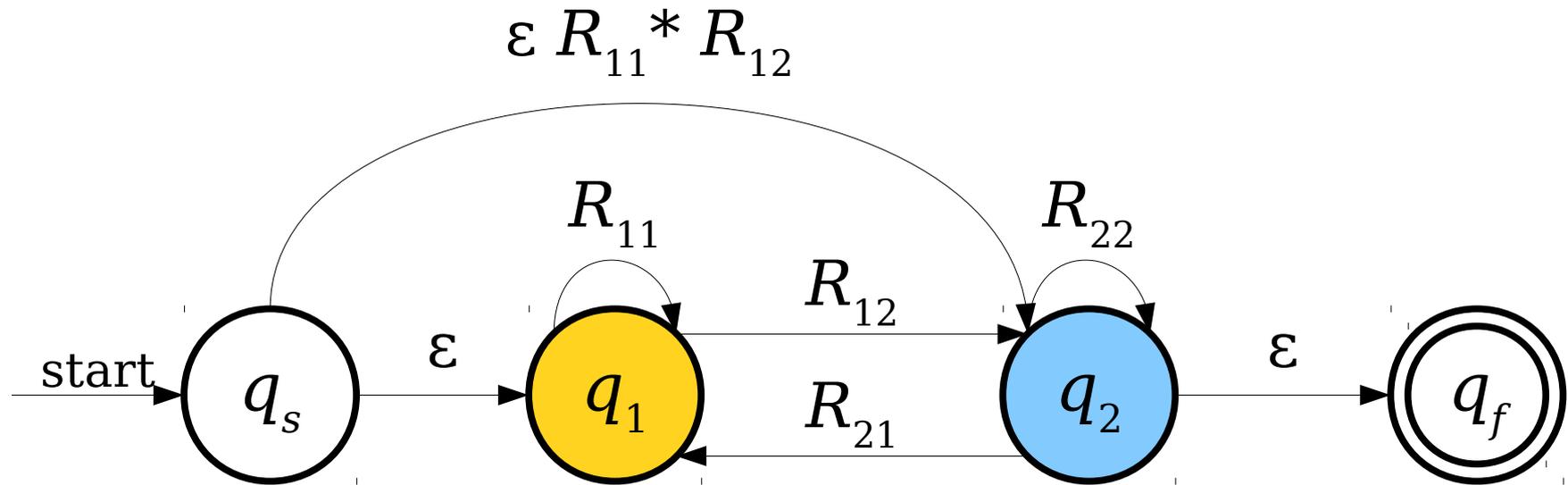


Note: We're using concatenation and Kleene closure in order to skip this state.

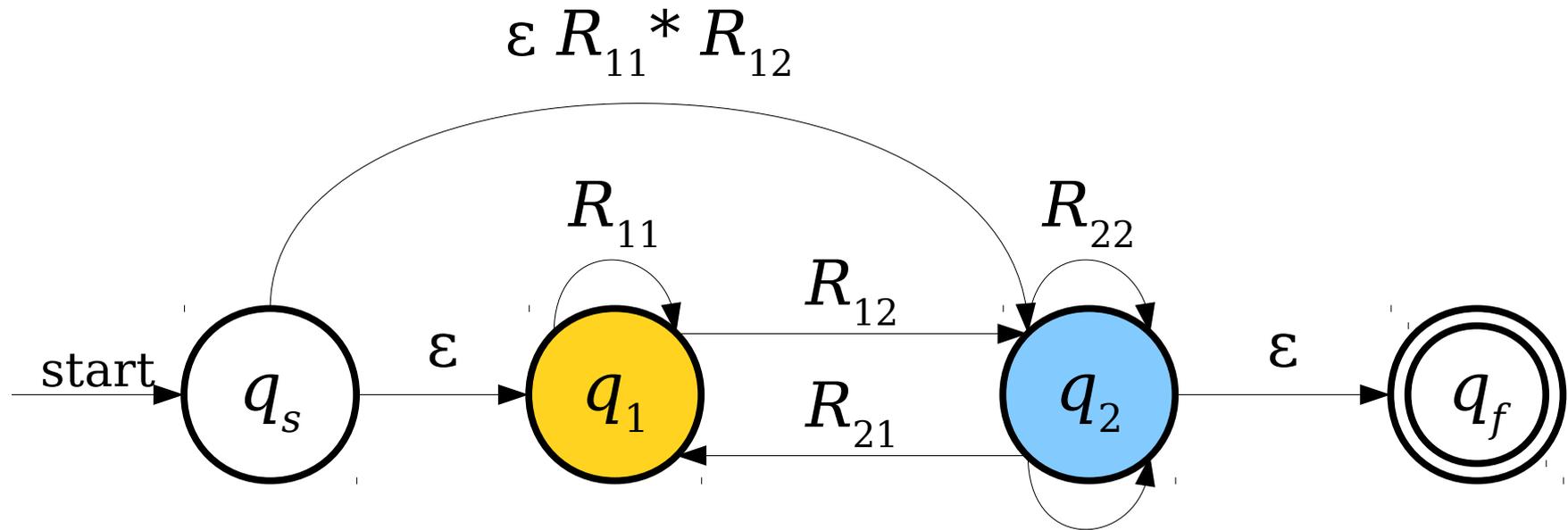
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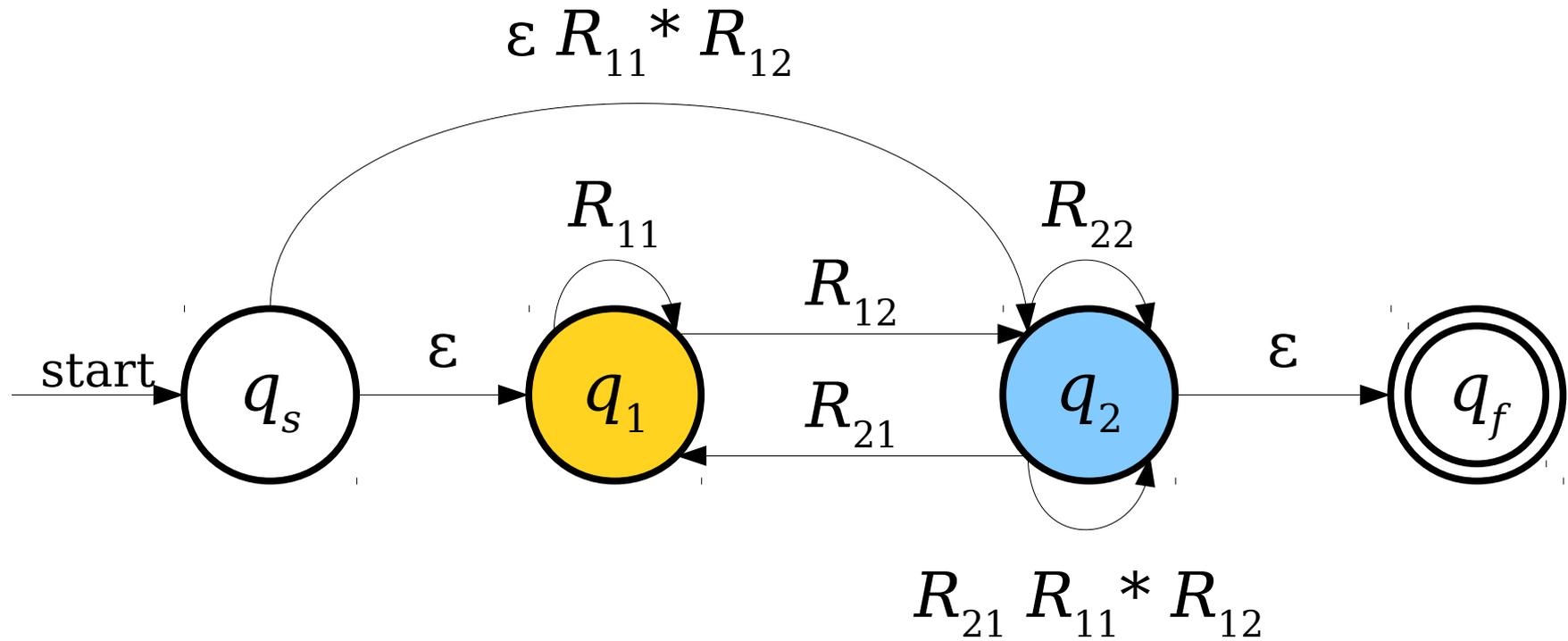
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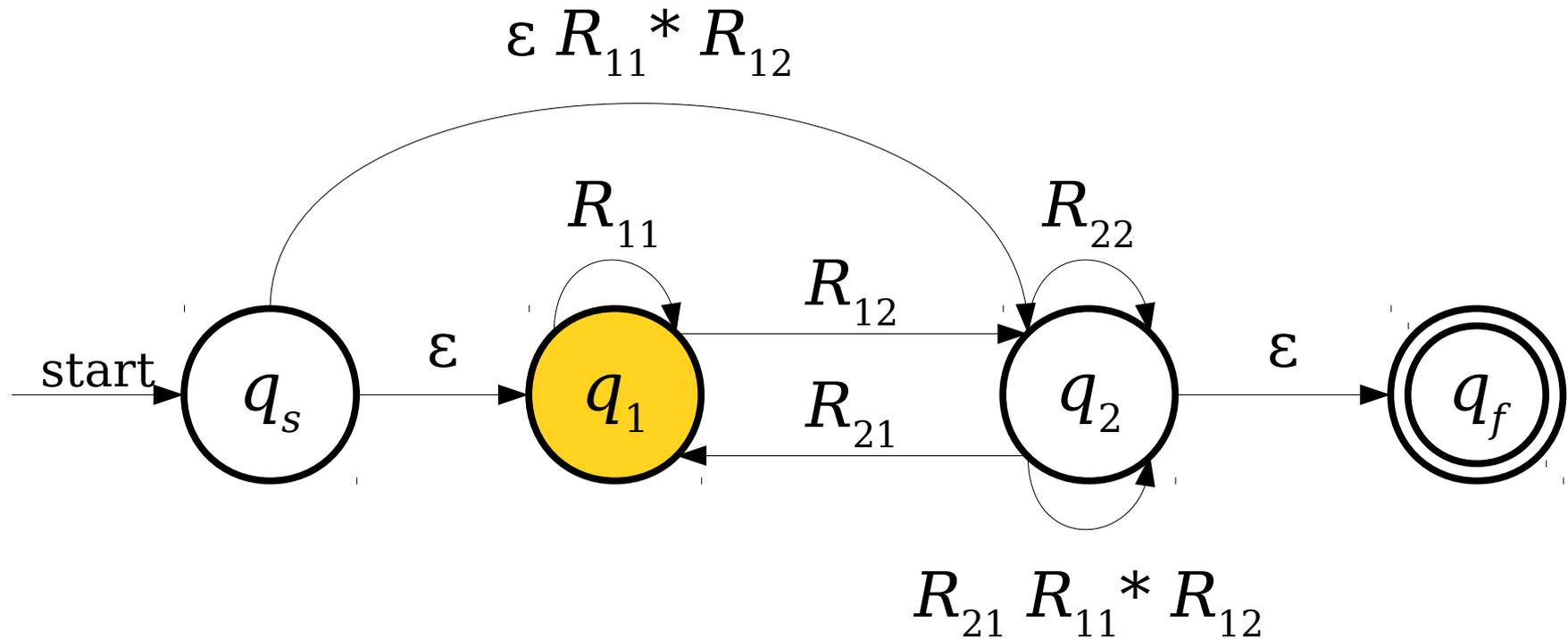
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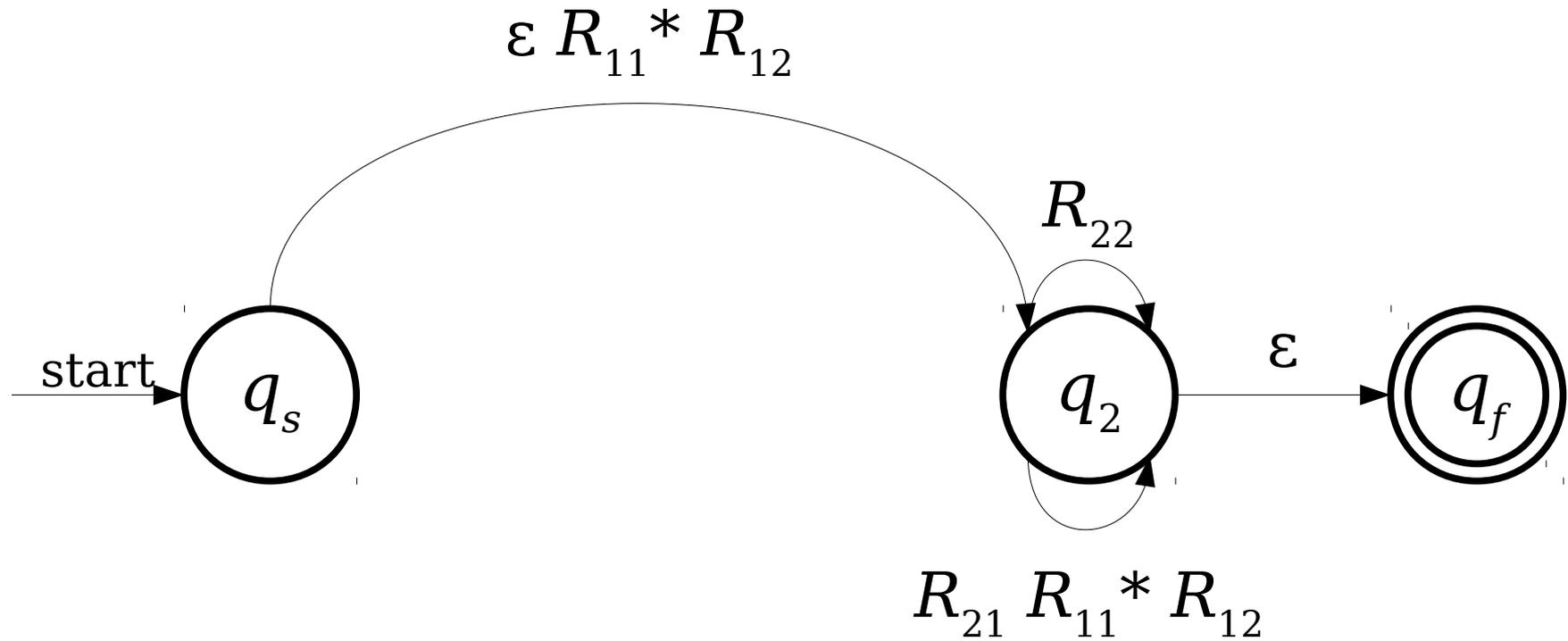
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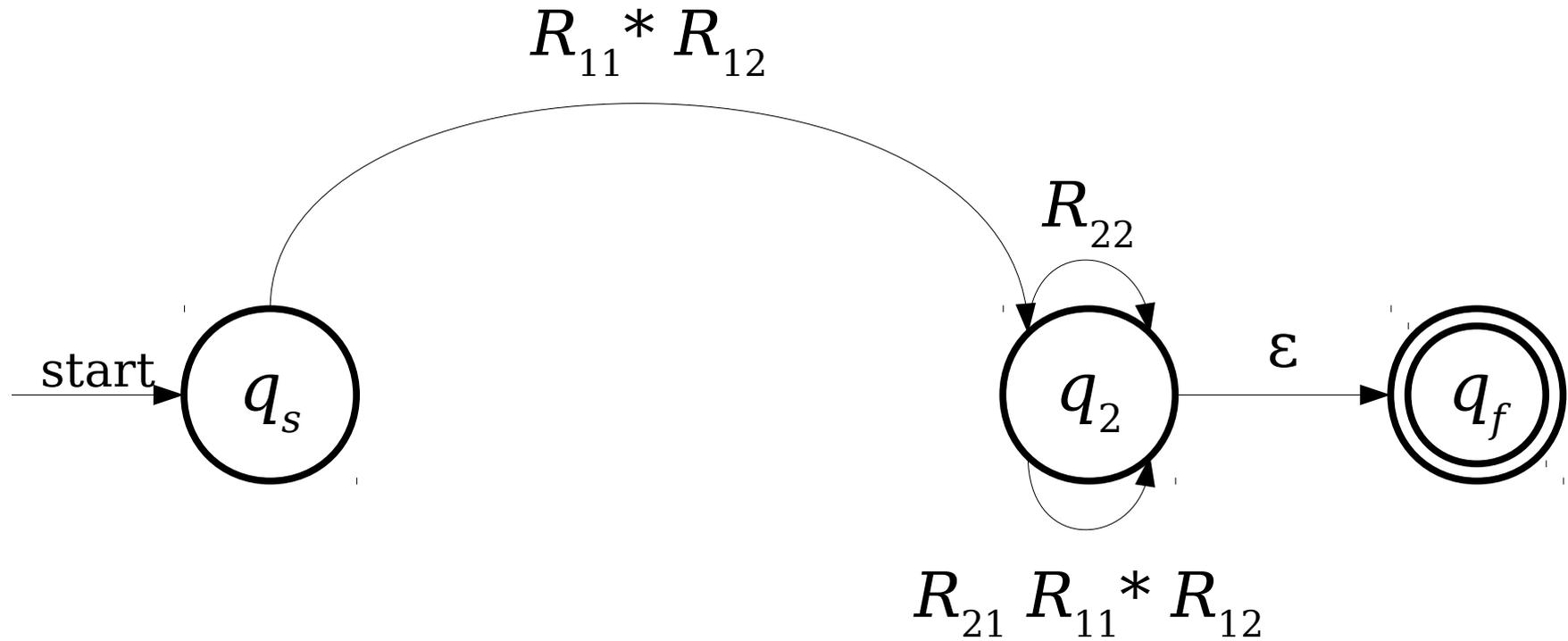
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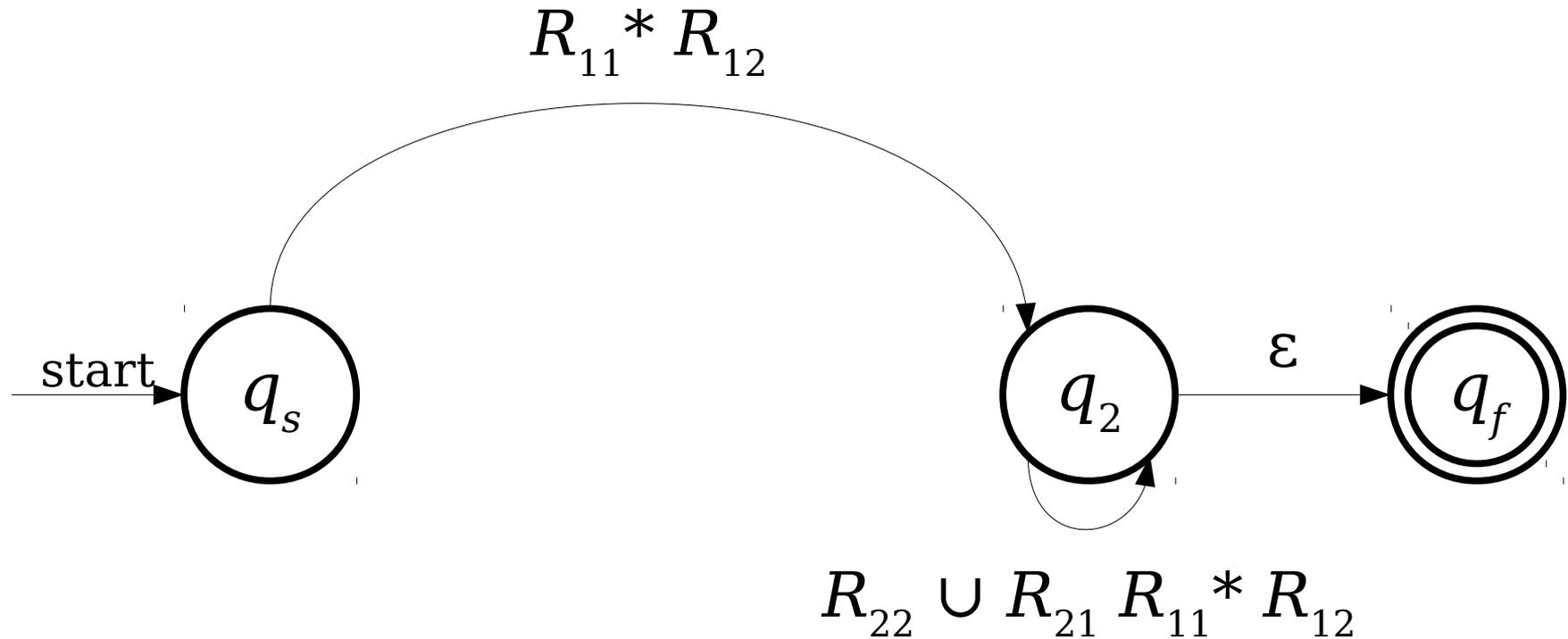
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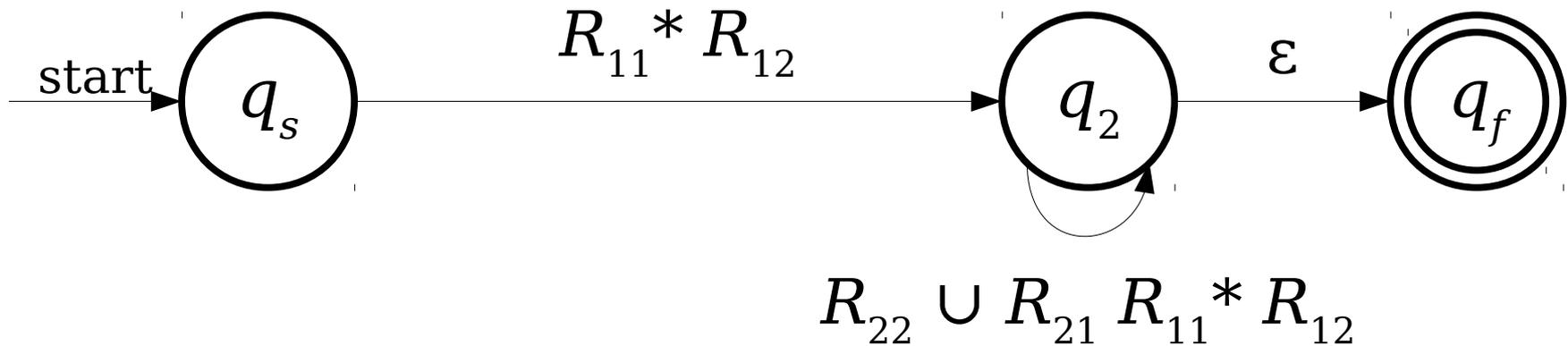


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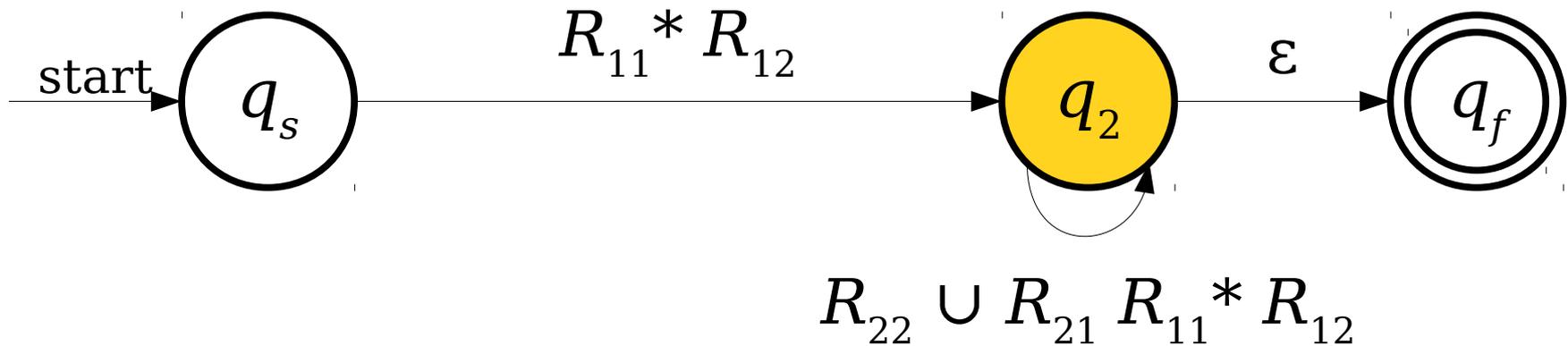


Note: We're using **union** to combine these transitions together.

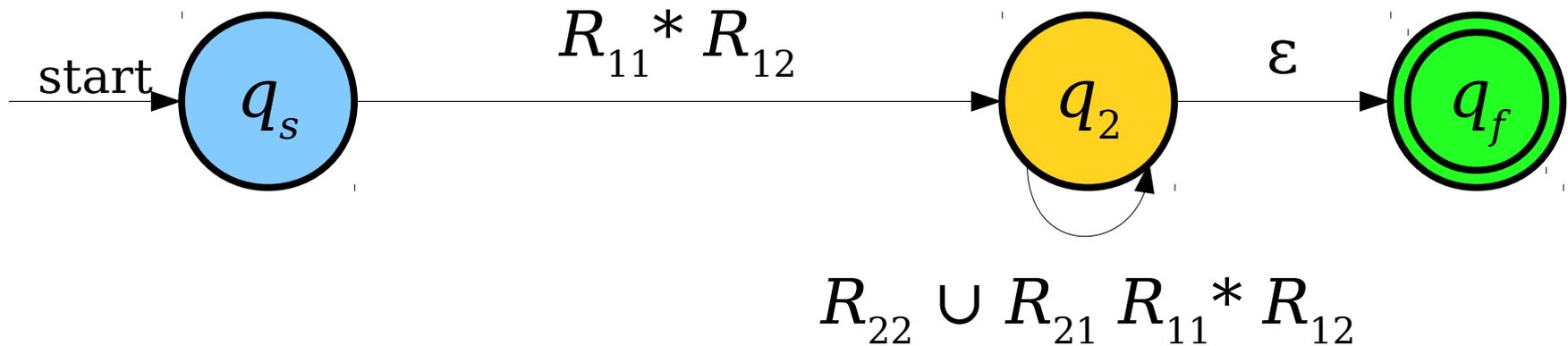
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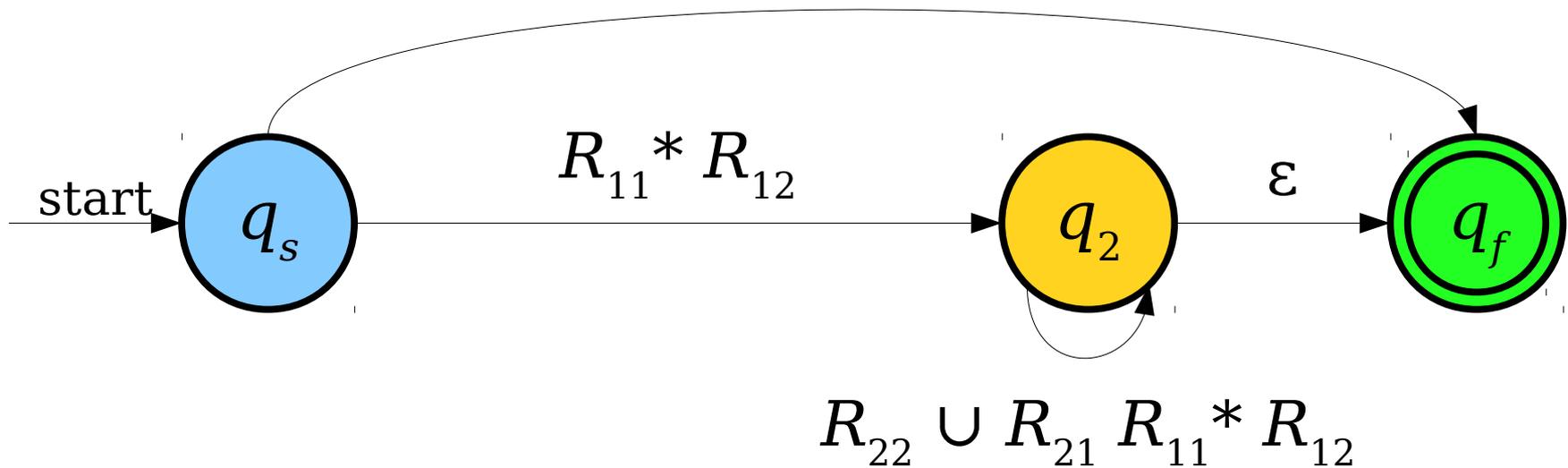
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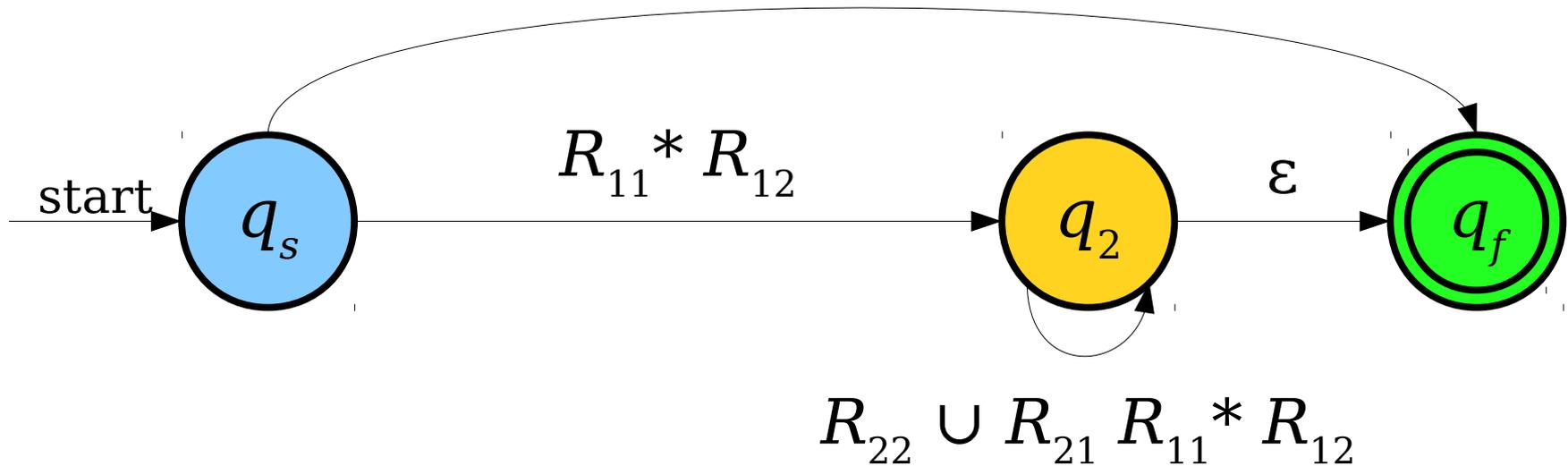


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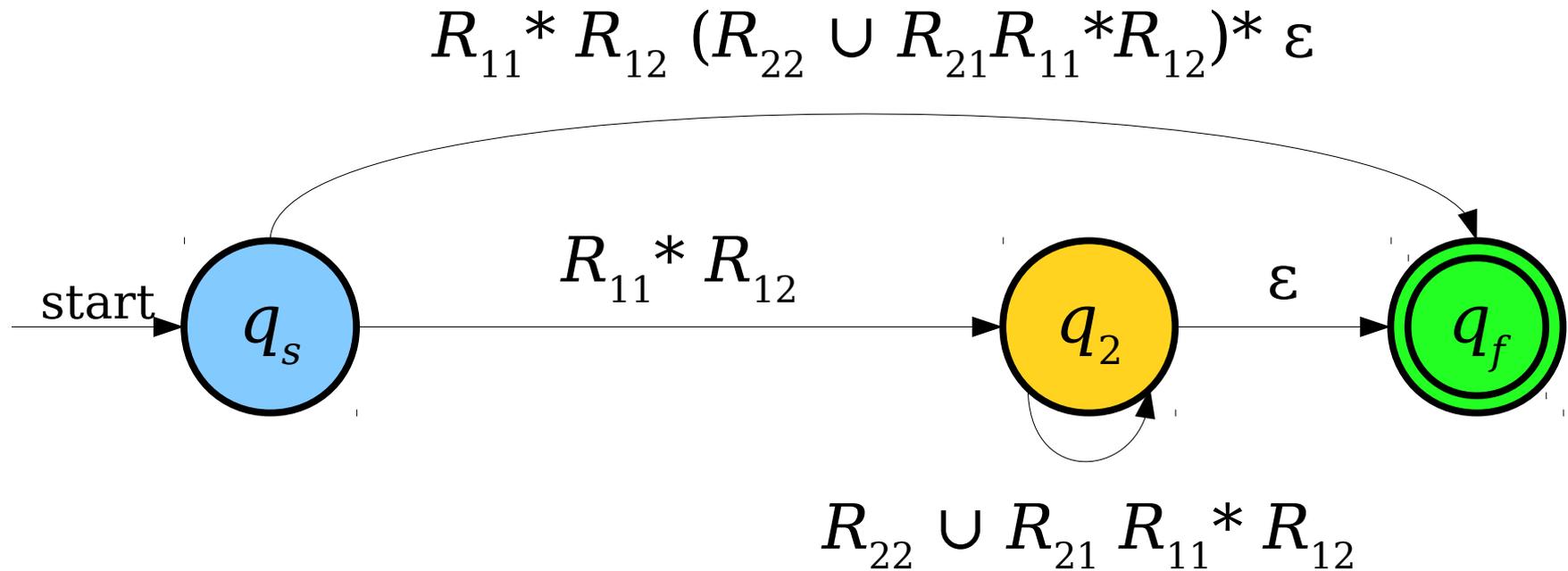


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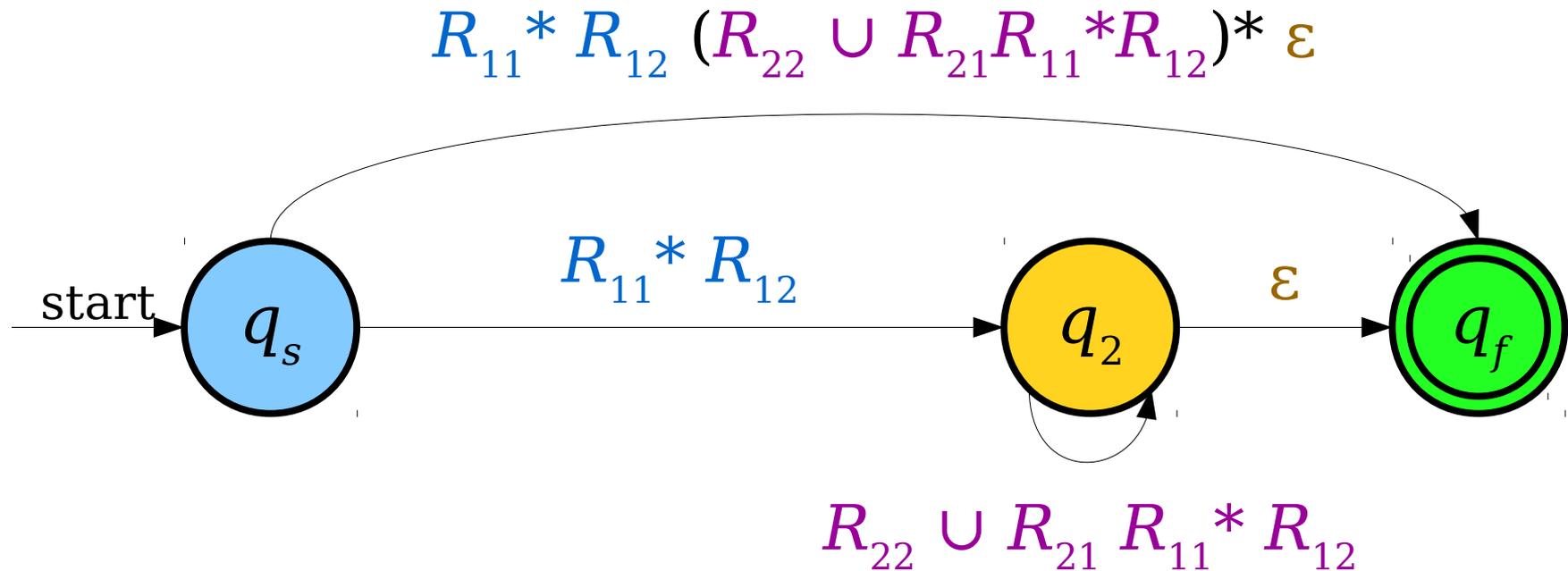
What should we put on this transition?



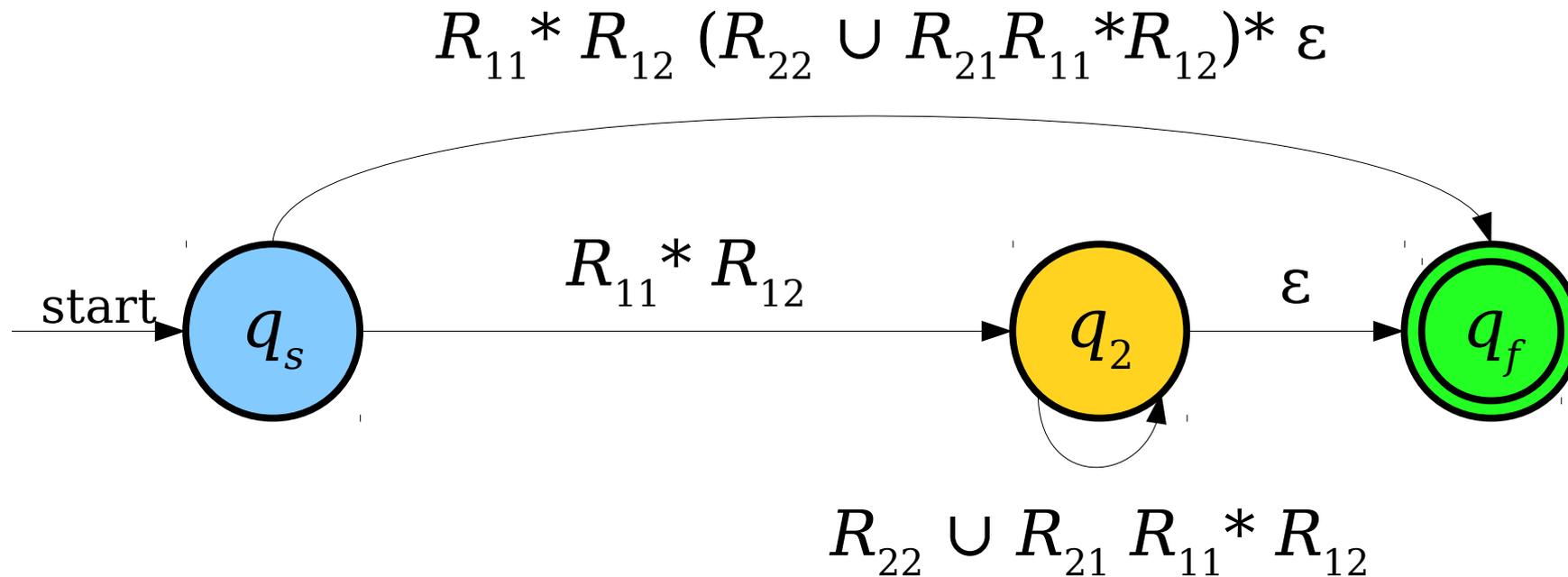
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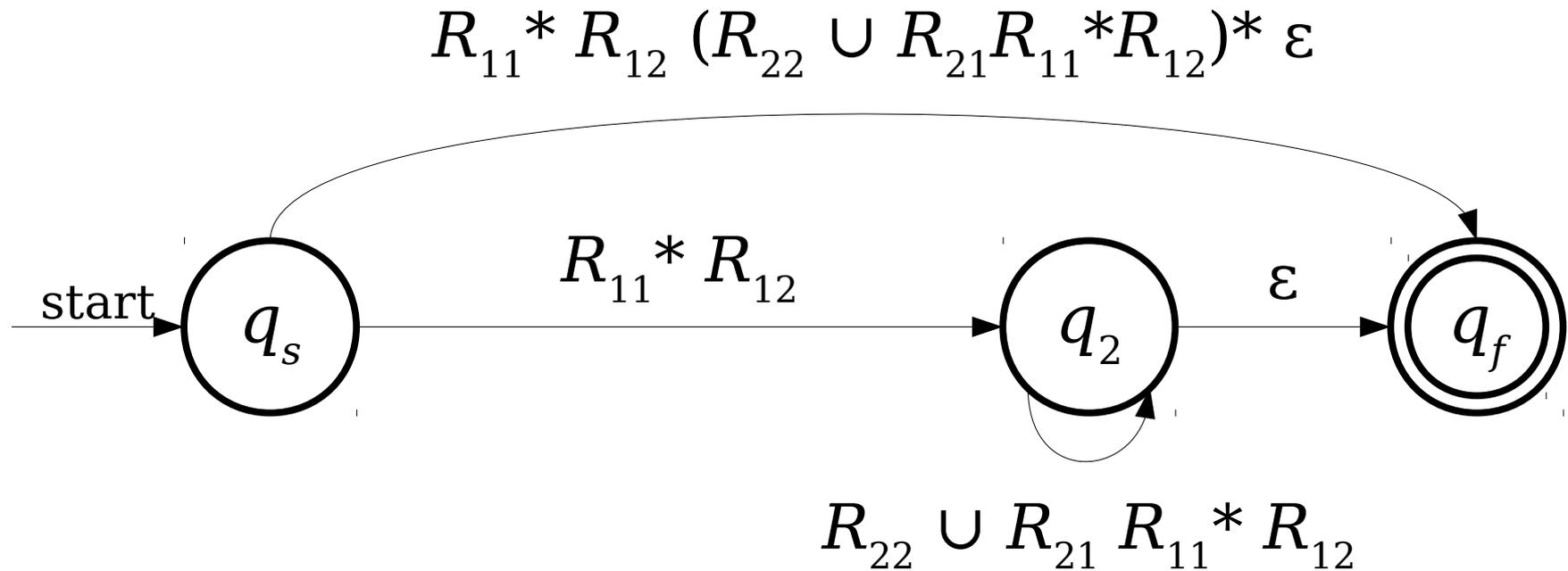
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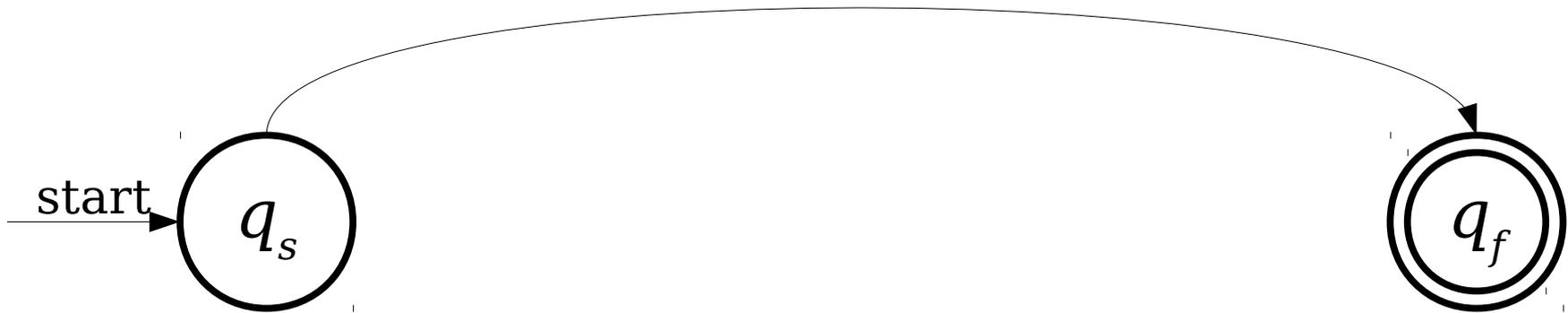


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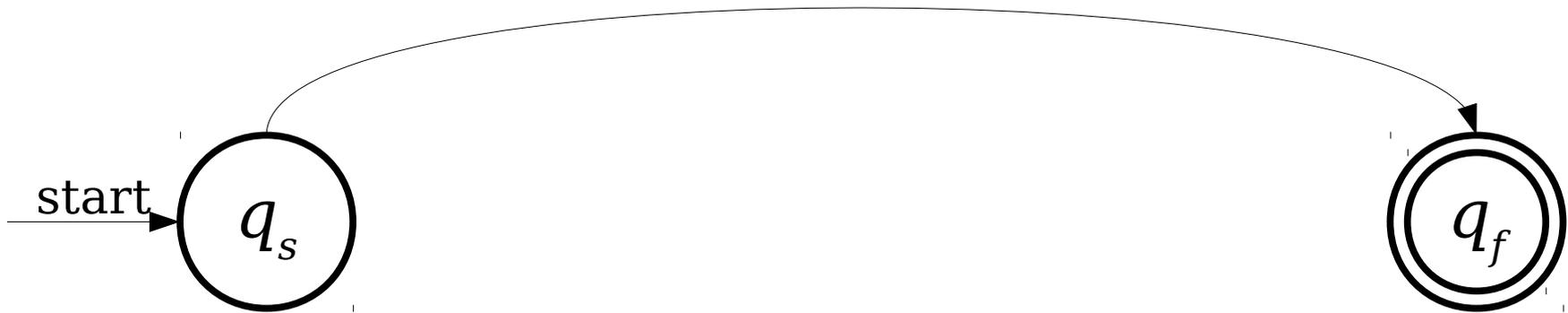
# From NFAs to Regular Expressions

$$R_{11}^* R_{12} (R_{22} \cup R_{21} R_{11}^* R_{12})^* \varepsilon$$

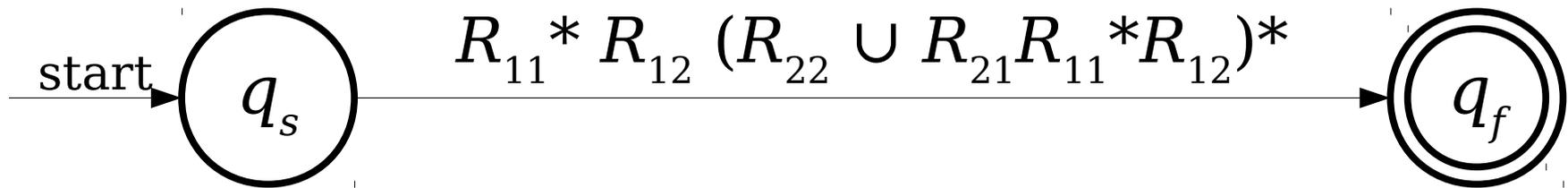


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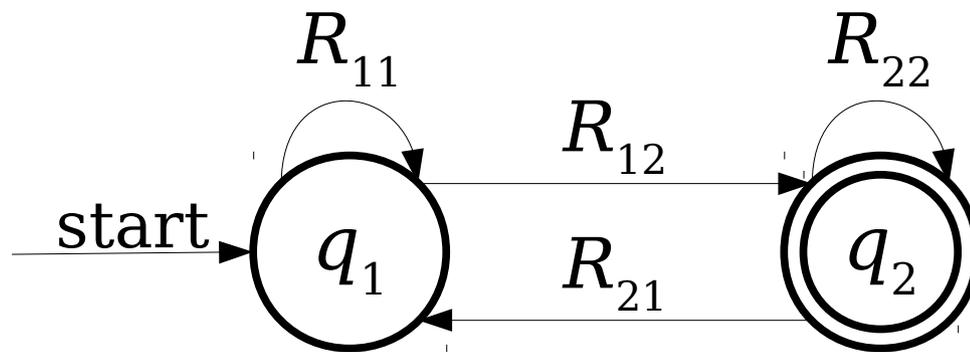
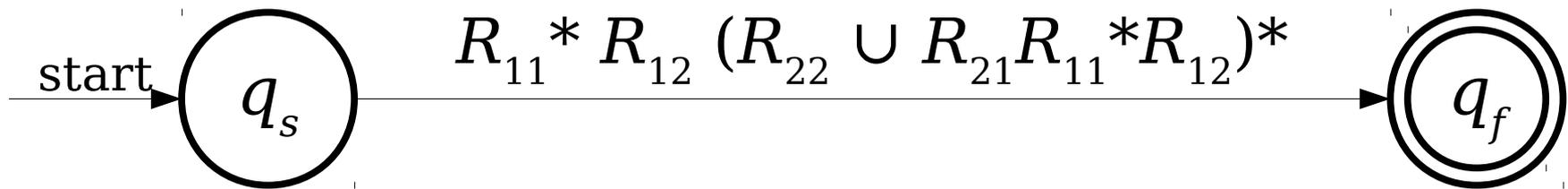
$$R_{11}^* R_{12} (R_{22} \cup R_{21} R_{11}^* R_{12})^*$$



# From NFAs to Regular Expressions



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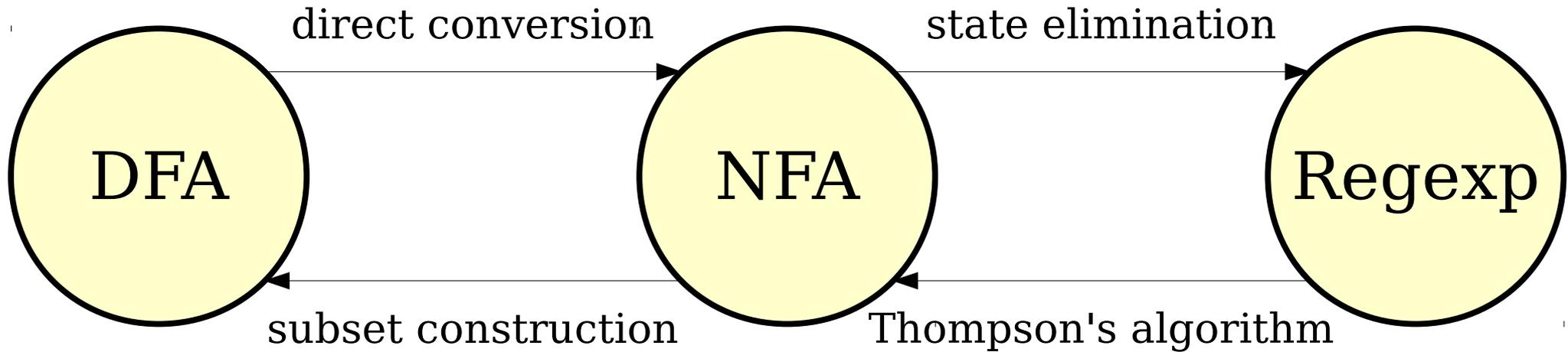
# The Construction at a Glance

- Start with an NFA  $N$  for the language  $L$ .
- Add a new start state  $q_s$  and accept state  $q_f$  to the NFA.
  - Add an  $\varepsilon$ -transition from  $q_s$  to the old start state of  $N$ .
  - Add  $\varepsilon$ -transitions from each accepting state of  $N$  to  $q_f$ , then mark them as not accepting.
- Repeatedly remove states other than  $q_s$  and  $q_f$  from the NFA by “shortcutting” them until only two states remain:  $q_s$  and  $q_f$ .
- The transition from  $q_s$  to  $q_f$  is then a regular expression for the NFA.

# Eliminating a State

- To eliminate a state  $q$  from the automaton, do the following for each pair of states  $q_0$  and  $q_1$ , where there's a transition from  $q_0$  into  $q$  and a transition from  $q$  into  $q_1$ :
  - Let  $R_{in}$  be the regex on the transition from  $q_0$  to  $q$ .
  - Let  $R_{out}$  be the regex on the transition from  $q$  to  $q_1$ .
  - If there is a regular expression  $R_{stay}$  on a transition from  $q$  to itself, add a new transition from  $q_0$  to  $q_1$  labeled  $((R_{in})(R_{stay})^*(R_{out}))$ .
  - If there isn't, add a new transition from  $q_0$  to  $q_1$  labeled  $((R_{in})(R_{out}))$
- If a pair of states has multiple transitions between them labeled  $R_1, R_2, \dots, R_k$ , replace them with a single transition labeled  $R_1 \cup R_2 \cup \dots \cup R_k$ .

# Our Transformations



**Theorem:** The following are all equivalent:

- $L$  is a regular language.
- There is a DFA  $D$  such that  $\mathcal{L}(D) = L$ .
- There is an NFA  $N$  such that  $\mathcal{L}(N) = L$ .
- There is a regular expression  $R$  such that  $\mathcal{L}(R) = L$ .

# Why This Matters

- The equivalence of regular expressions and finite automata has practical relevance.
  - Tools like `grep` and `flex` that use regular expressions capture all the power available via DFAs and NFAs.
- This also is hugely theoretically significant: the regular languages can be assembled “from scratch” using a small number of operations!

# Next Time

- ***Applications of Regular Languages***
  - Answering “so what?”
- ***Intuiting Regular Languages***
  - What makes a language regular?
- ***The Myhill-Nerode Theorem***
  - The limits of regular languages.