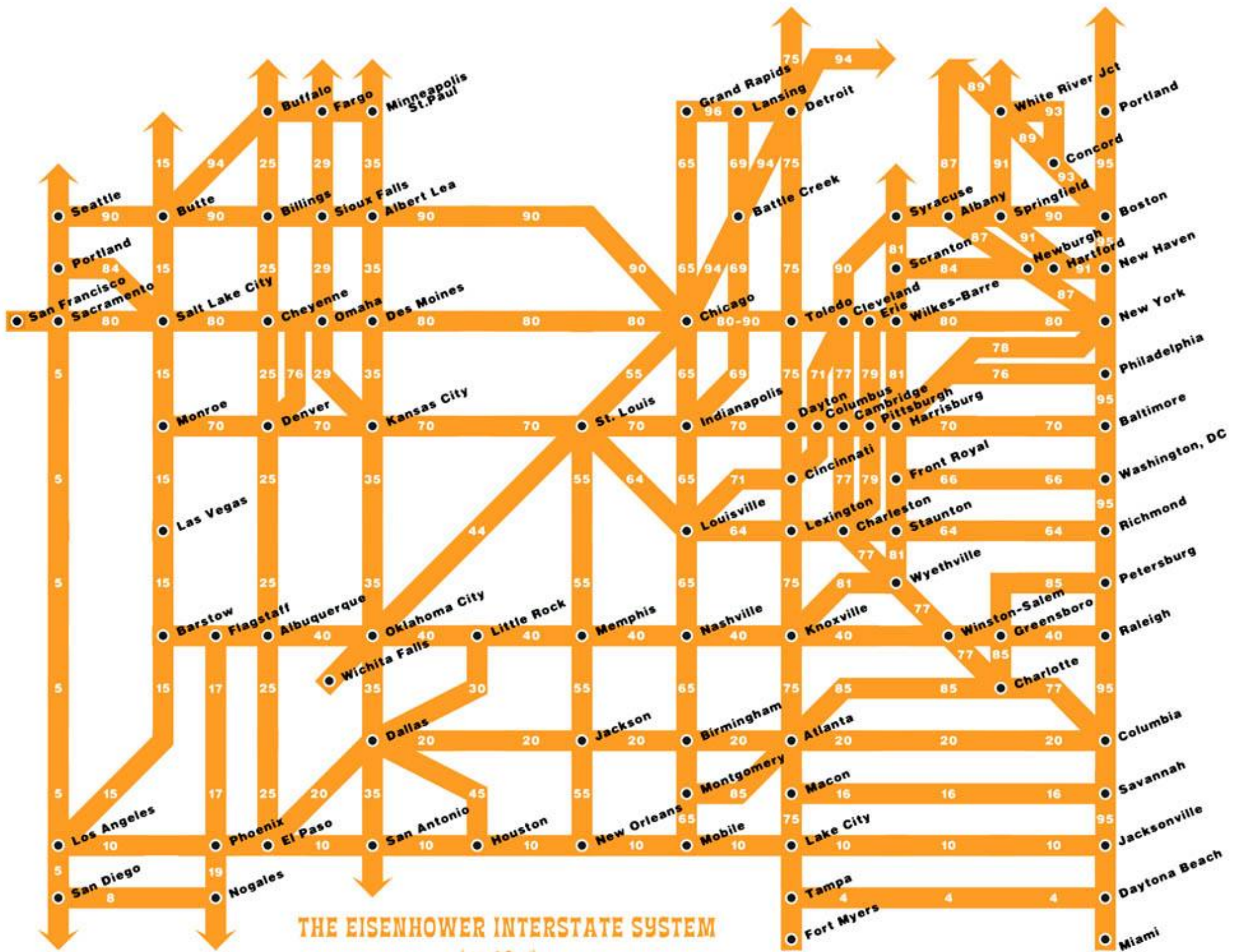


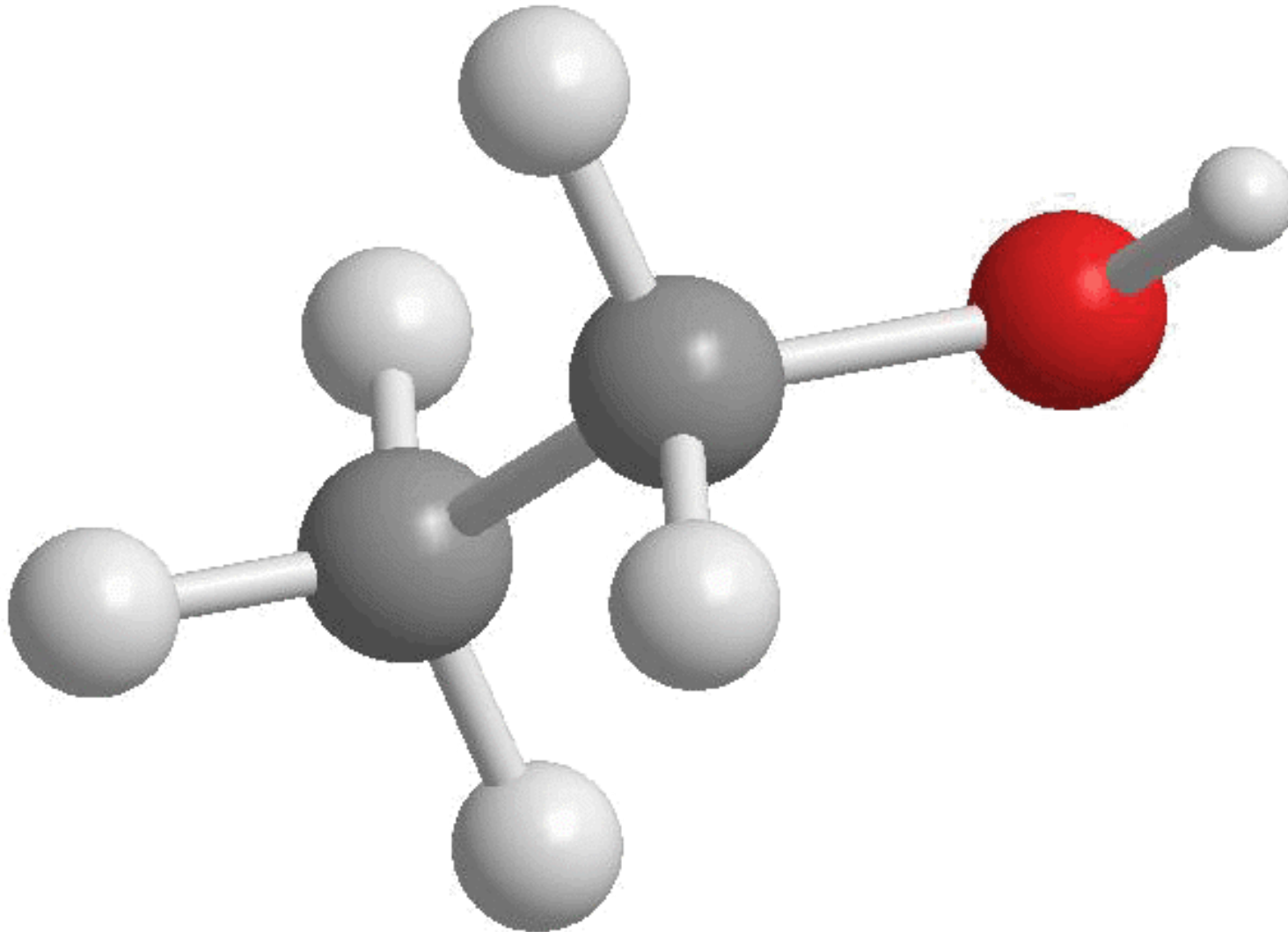
# Graph Theory

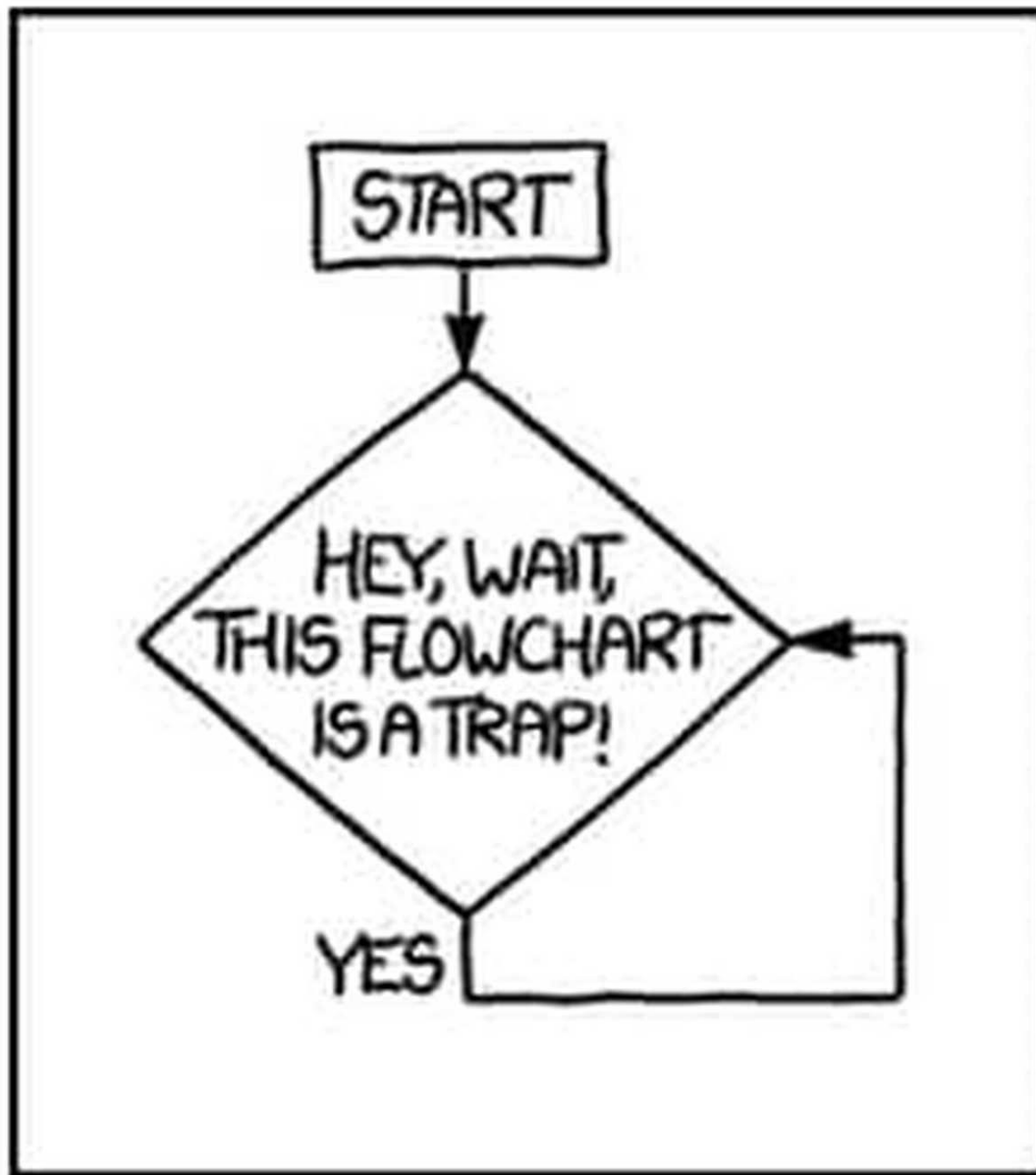


**THE EISENHOWER INTERSTATE SYSTEM**

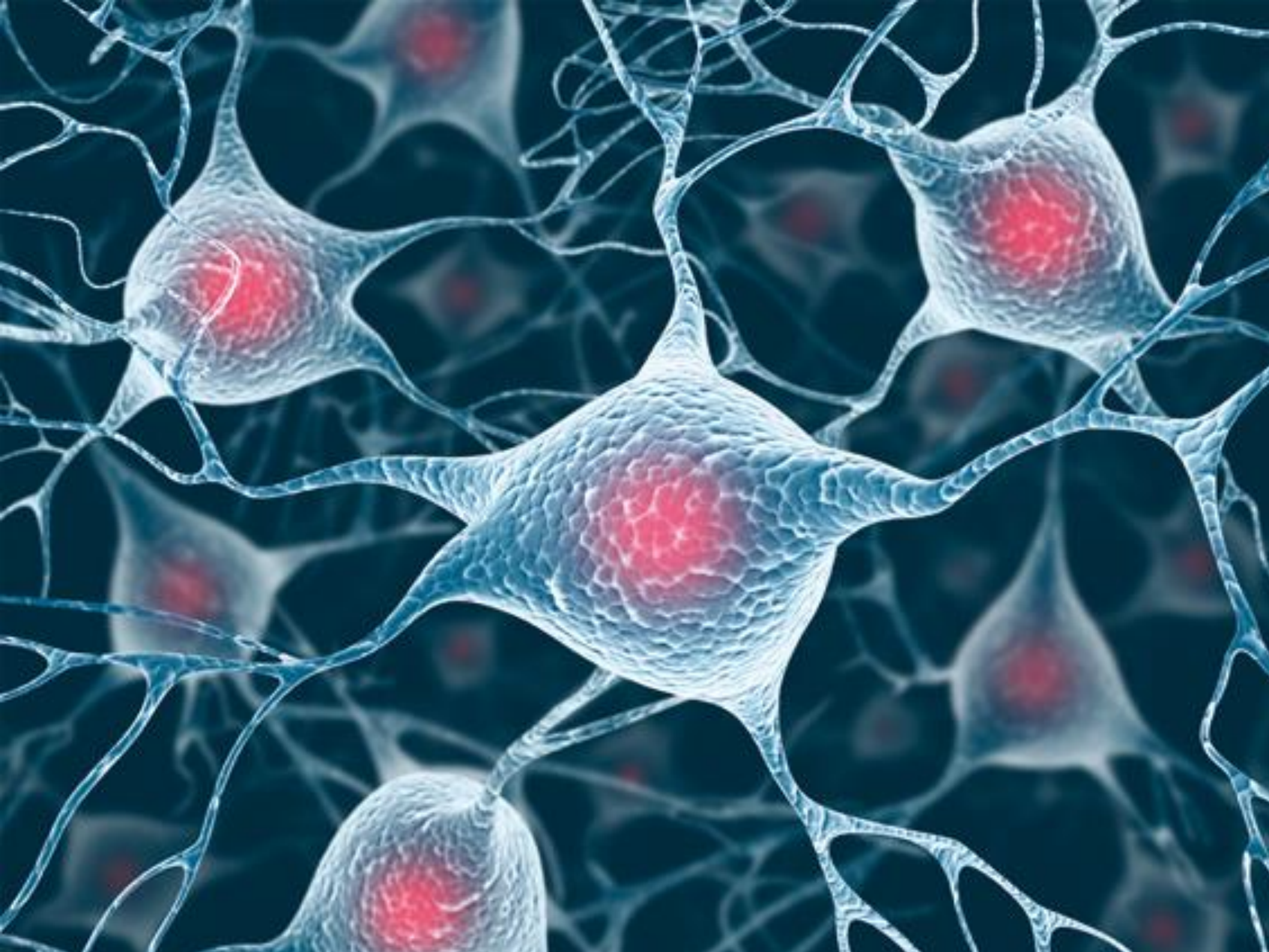
(simplified)

# Chemical Bonds









**facebook**®





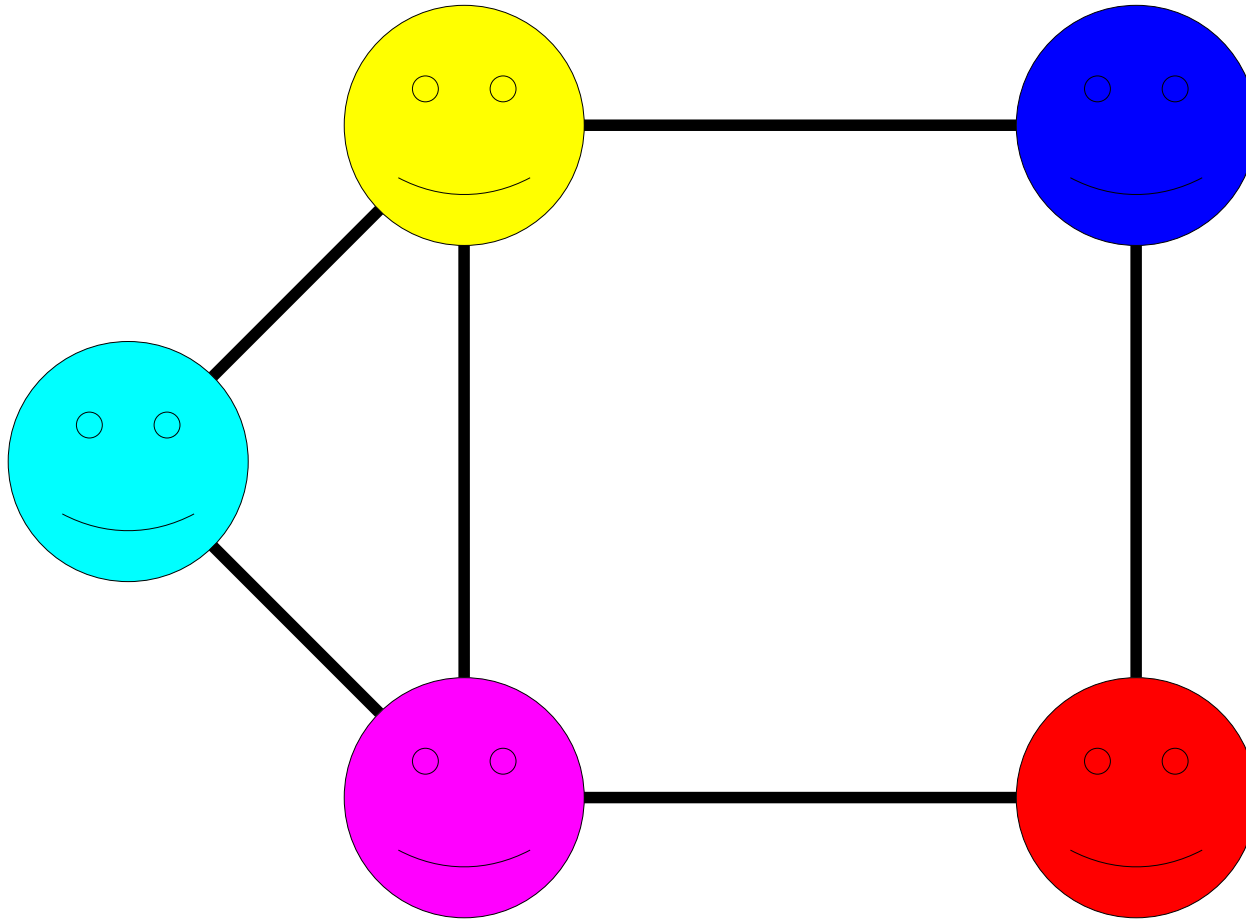
# What's in Common

Each of these structures consists of a collection of objects and links between those objects.

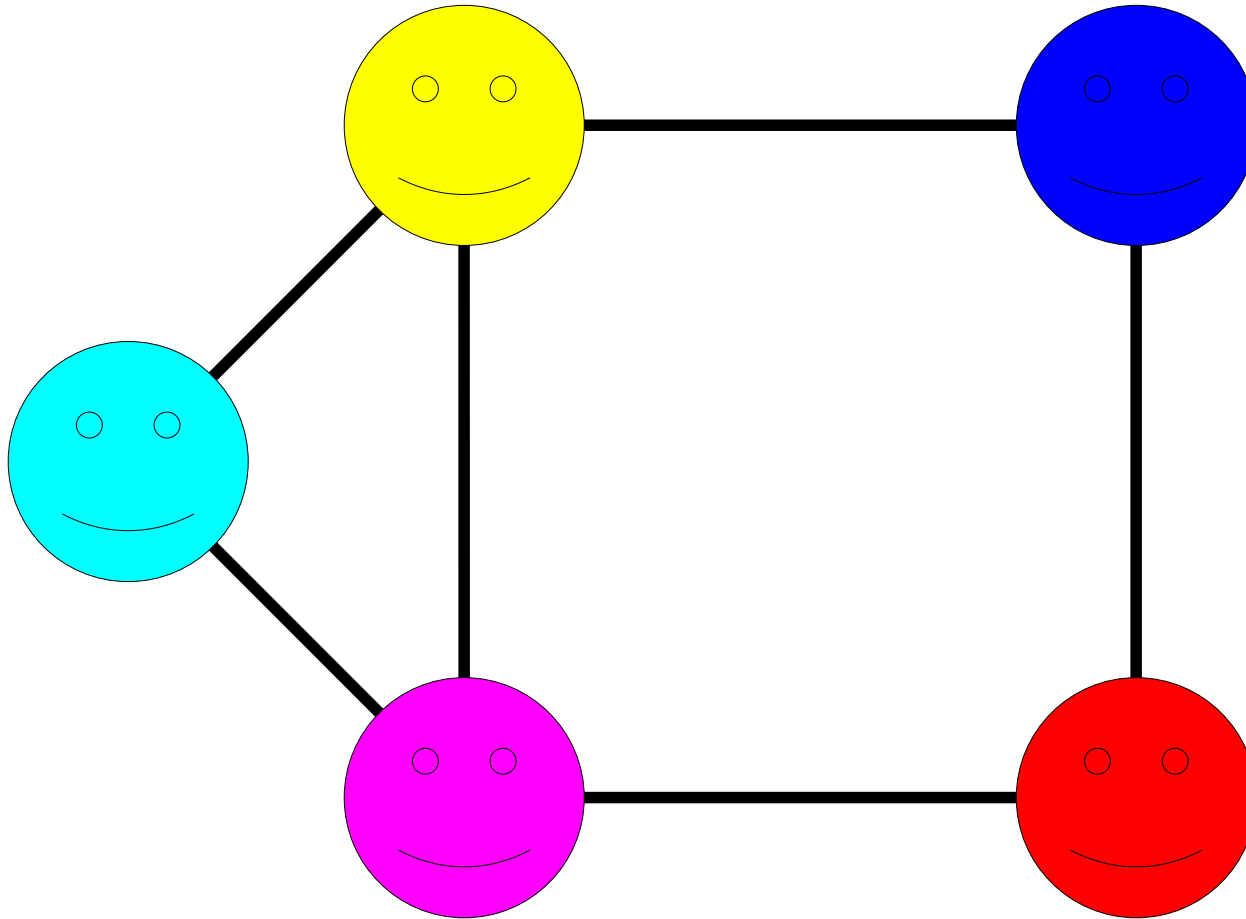
***Goal:*** find a general framework for describing these objects and their properties.



A **graph** is a mathematical structure for representing relationships.

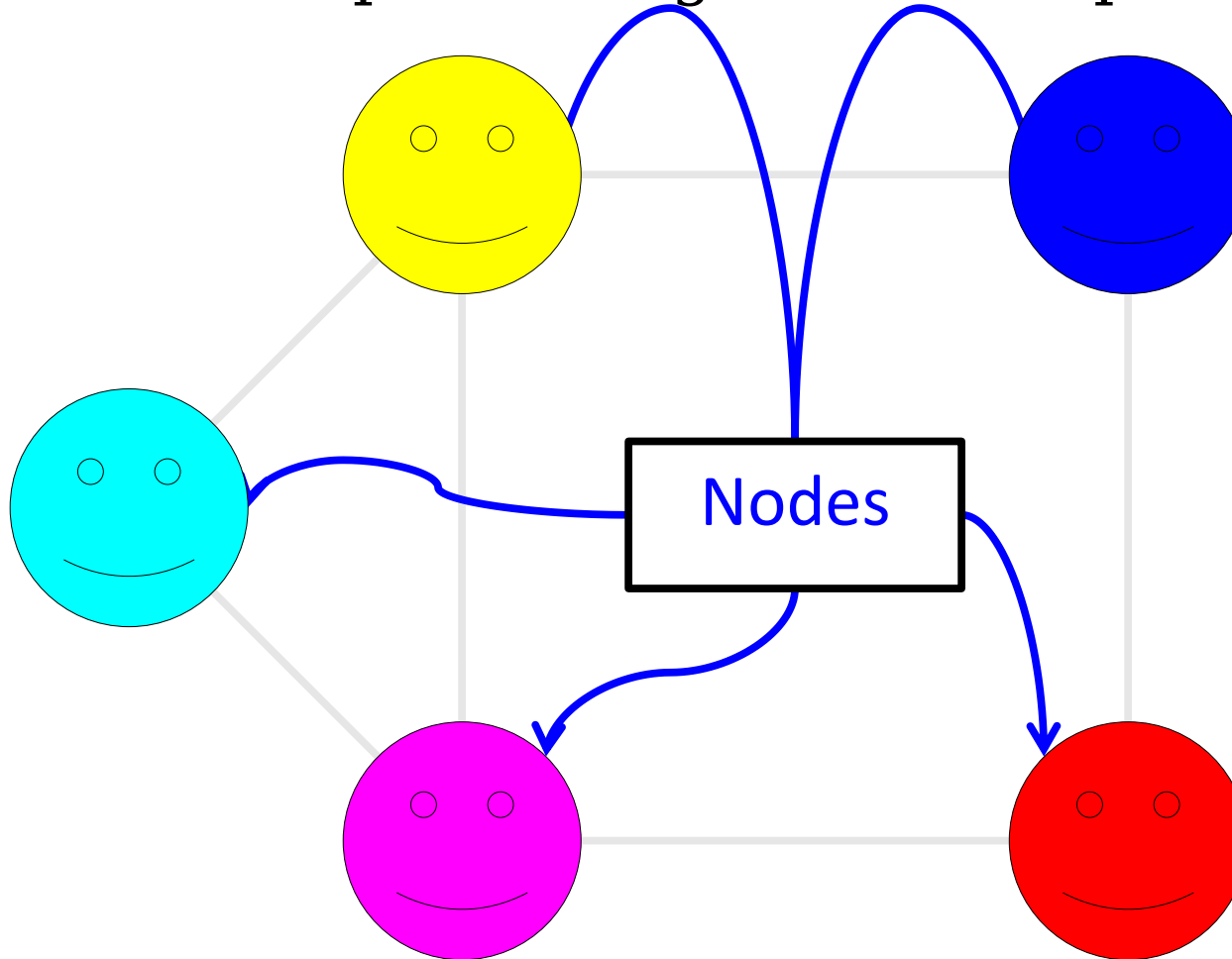


A **graph** is a mathematical structure for representing relationships.



A graph consists of a set of **nodes** (or **vertices**) connected by **edges** (or **arcs**)

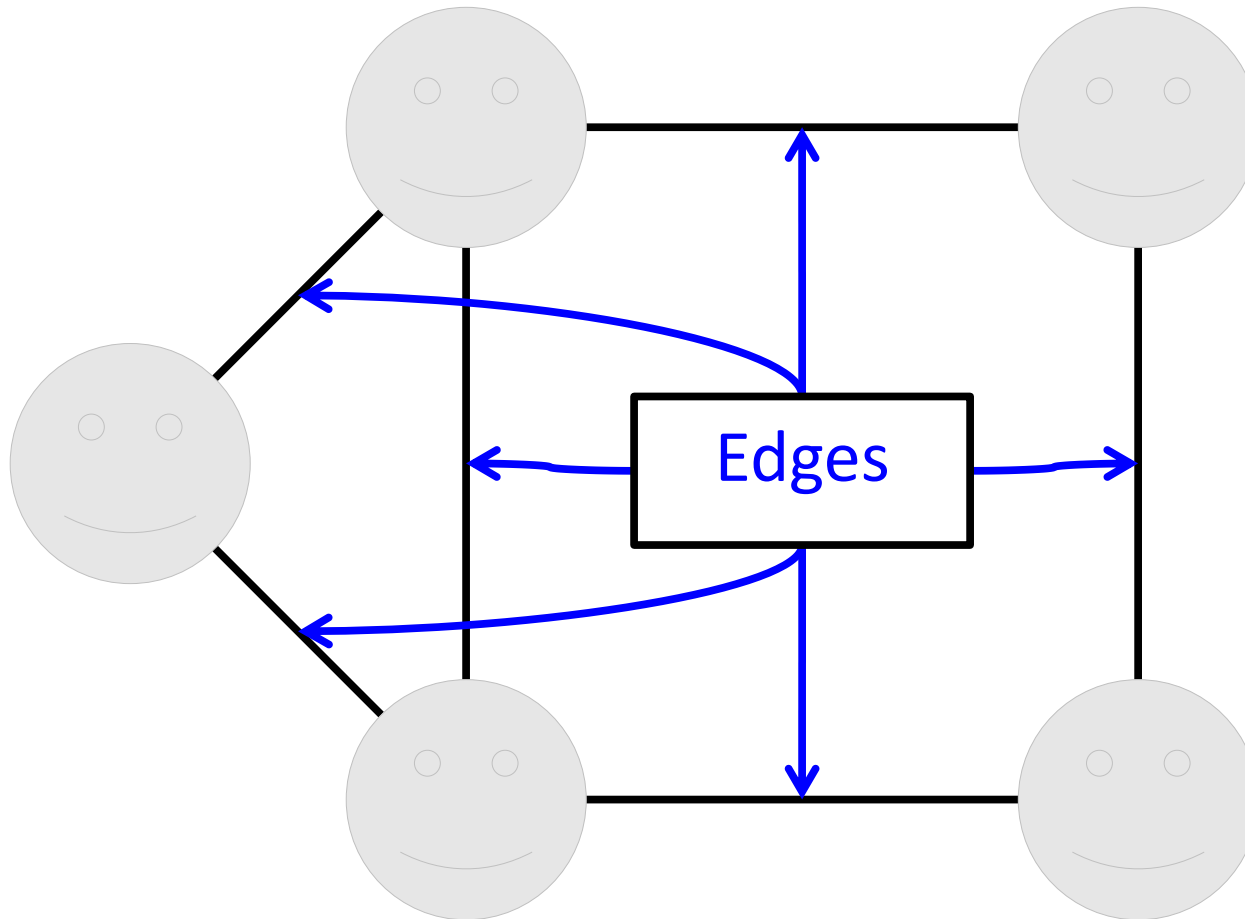
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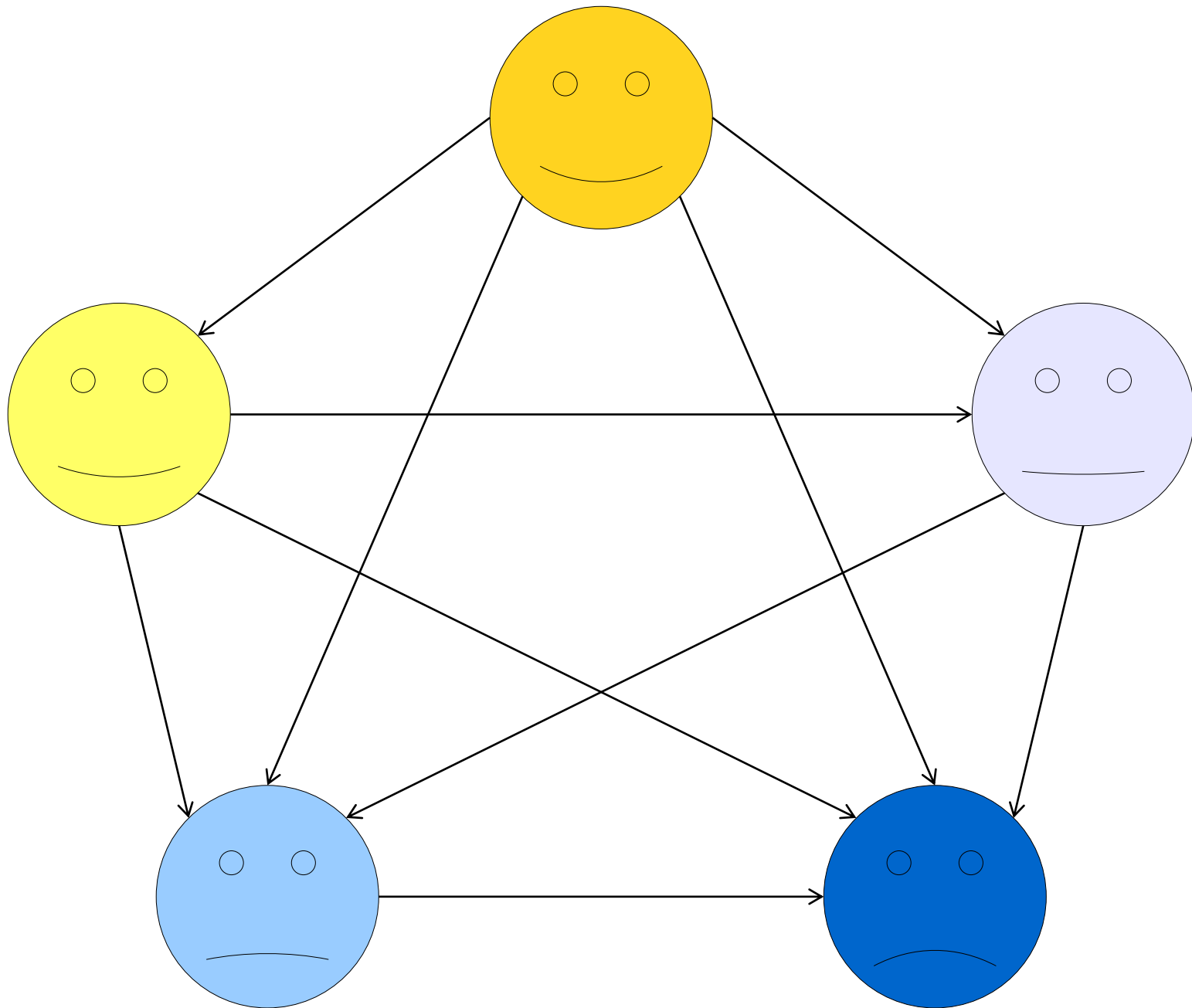


A **graph** is a mathematical structure for representing relationships.

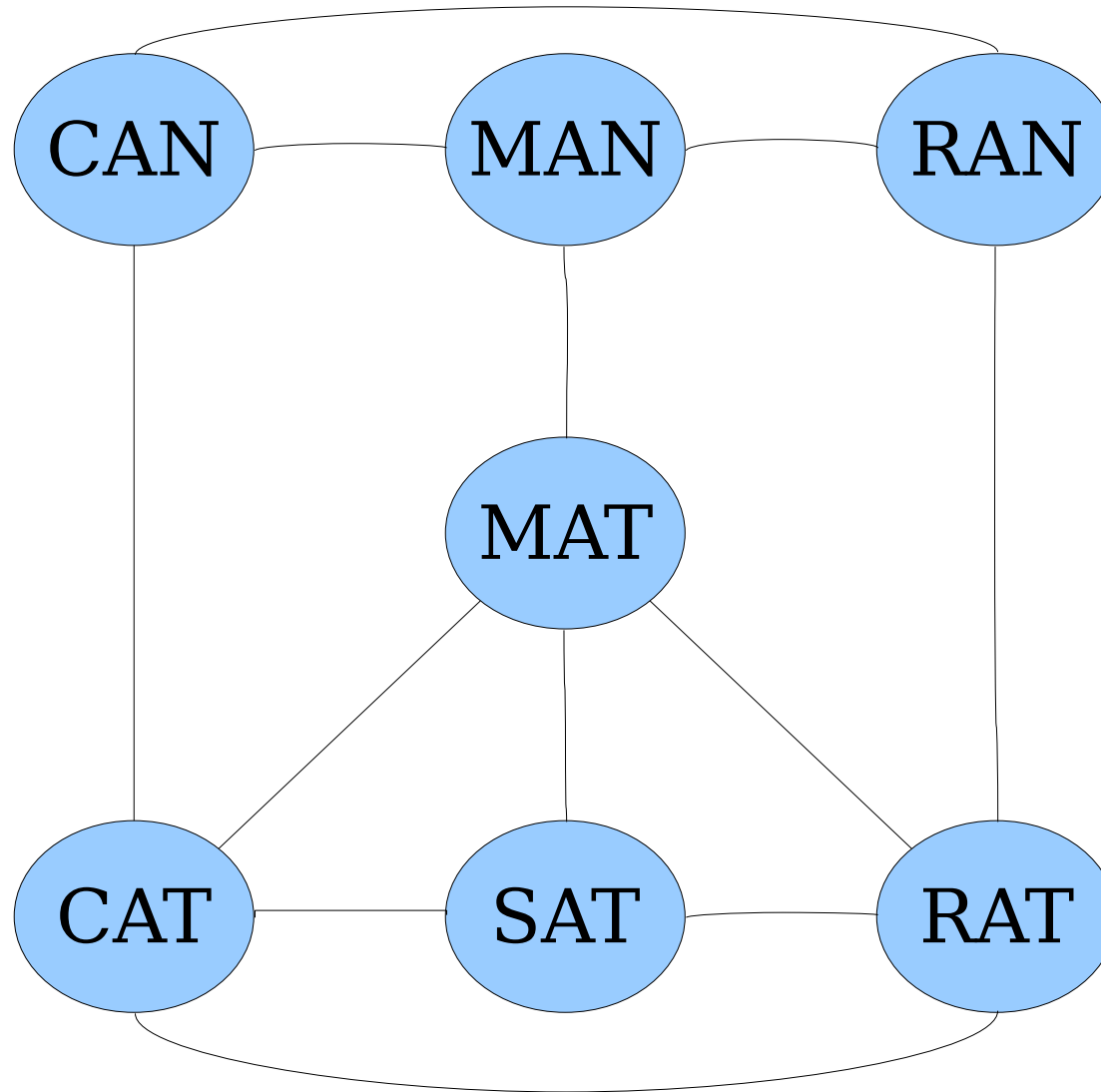


A graph consists of a set of **nodes** (or **vertices**) connected by **edges** (or **arcs**)

Some graphs are *directed*.



Some graphs are *undirected*.





Going forward, we're primarily going to focus on undirected graphs.

The term “graph” generally refers to undirected graphs with a finite number of nodes, unless specified otherwise.

# Formalizing Graphs

How might we define a graph mathematically?

We need to specify

- what the nodes in the graph are, and
- which edges are in the graph.

The nodes can be pretty much anything.

What about the edges?

# Formalizing Graphs

An **unordered pair** is a set  $\{a, b\}$  of two elements  $a \neq b$ . (Remember that sets are unordered).

$$\{0, 1\} = \{1, 0\}$$

An **undirected graph** is an ordered pair  $G = (V, E)$ , where

$V$  is a set of nodes, which can be anything, and

$E$  is a set of edges, which are unordered pairs of nodes drawn from  $V$ .

A **directed graph** is an ordered pair  $G = (V, E)$ , where

$V$  is a set of nodes, which can be anything, and

$E$  is a set of edges, which are *ordered* pairs of nodes drawn from  $V$ .



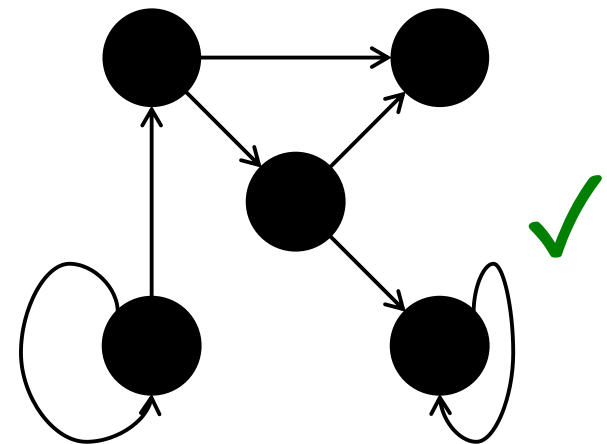
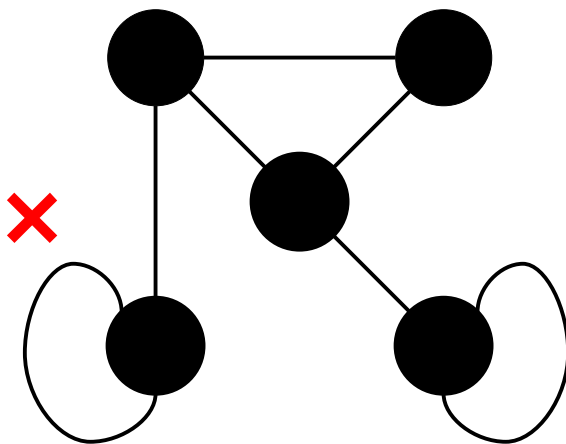
# Self-Loops

An edge from a node to itself is called a ***self-loop***.

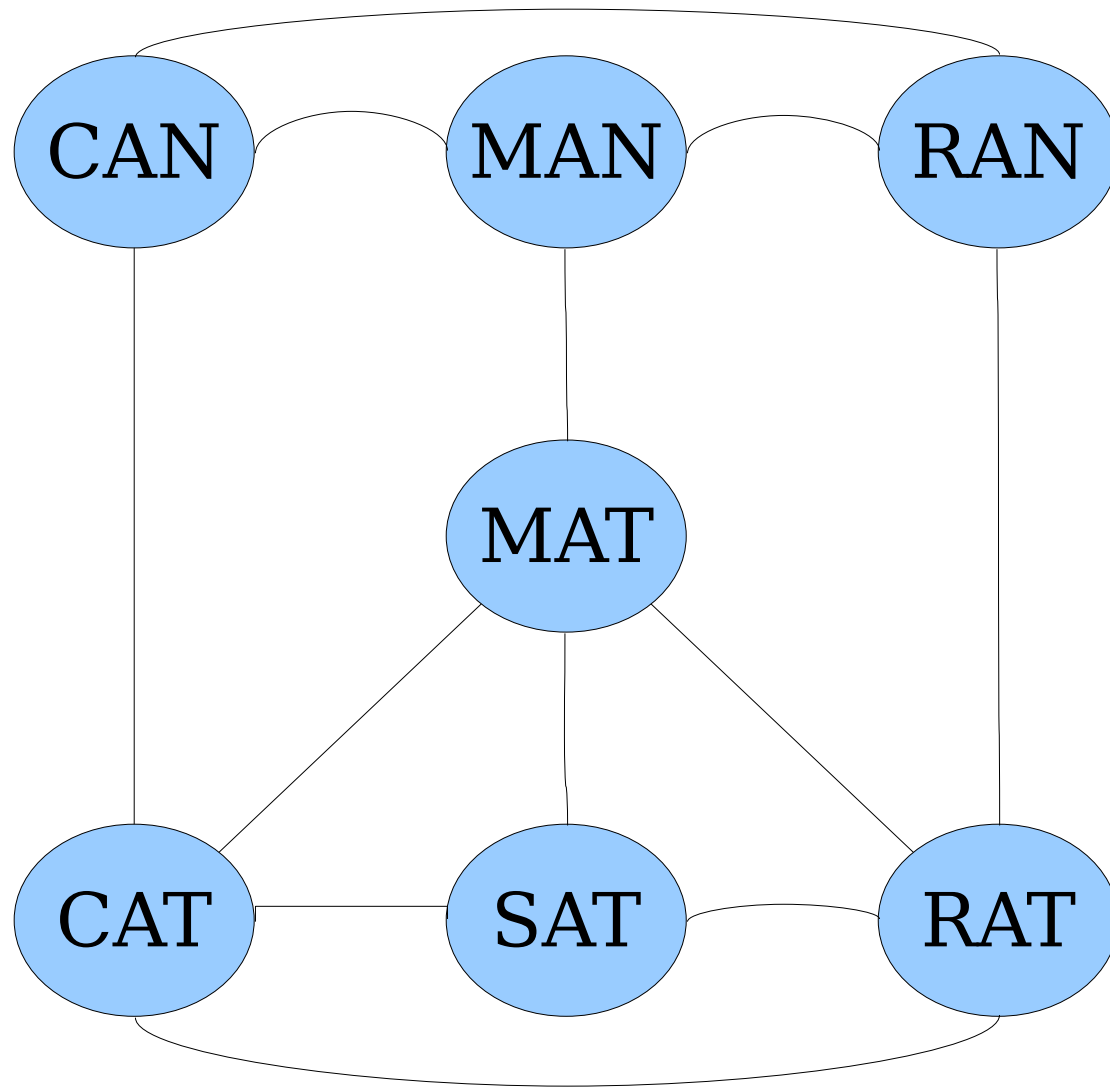
In undirected graphs, self-loops are generally not allowed.

Can you see how this follows from the definition?

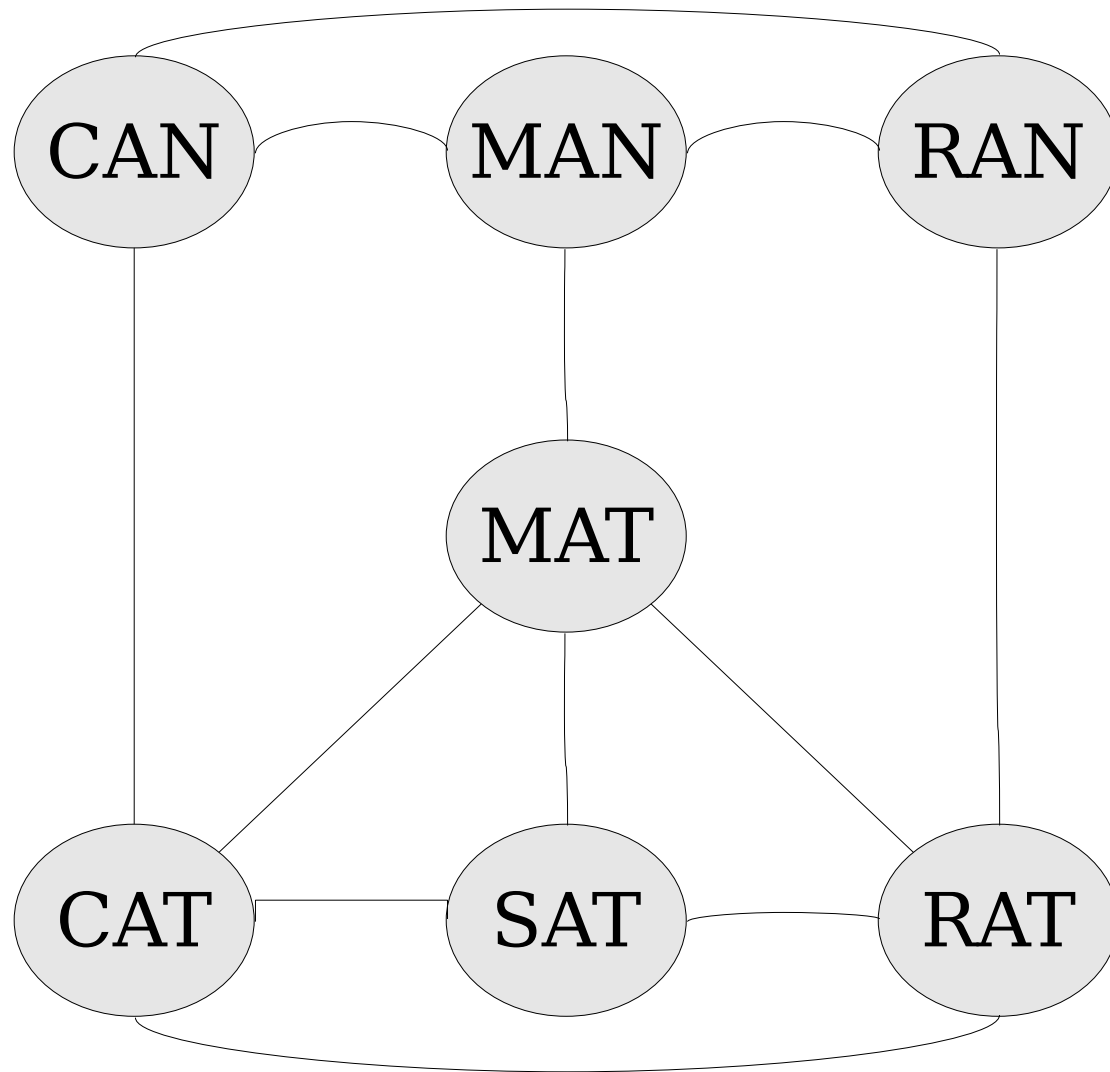
In directed graphs, self-loops are generally allowed unless specified otherwise.



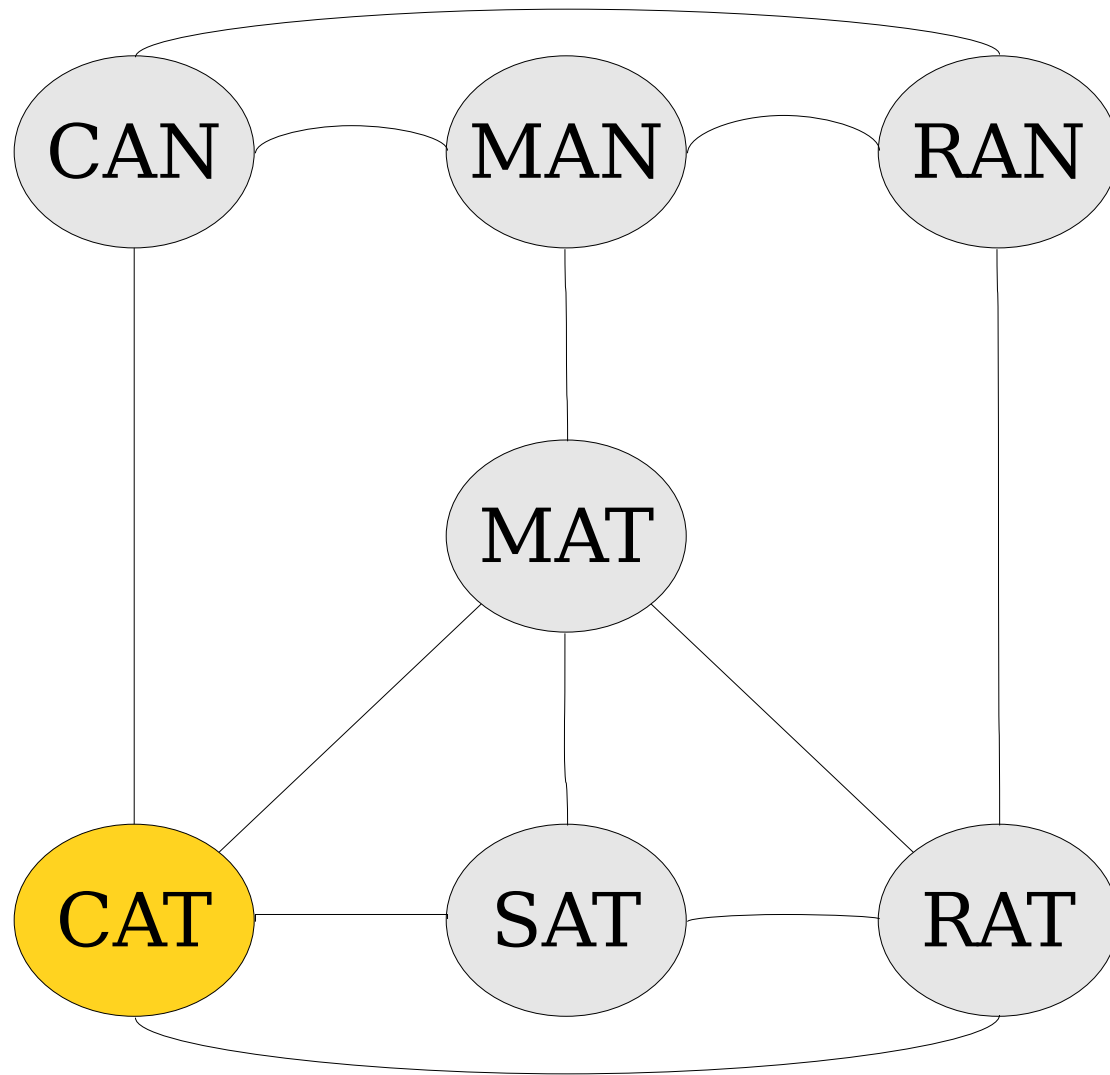
# Standard Graph Terminology



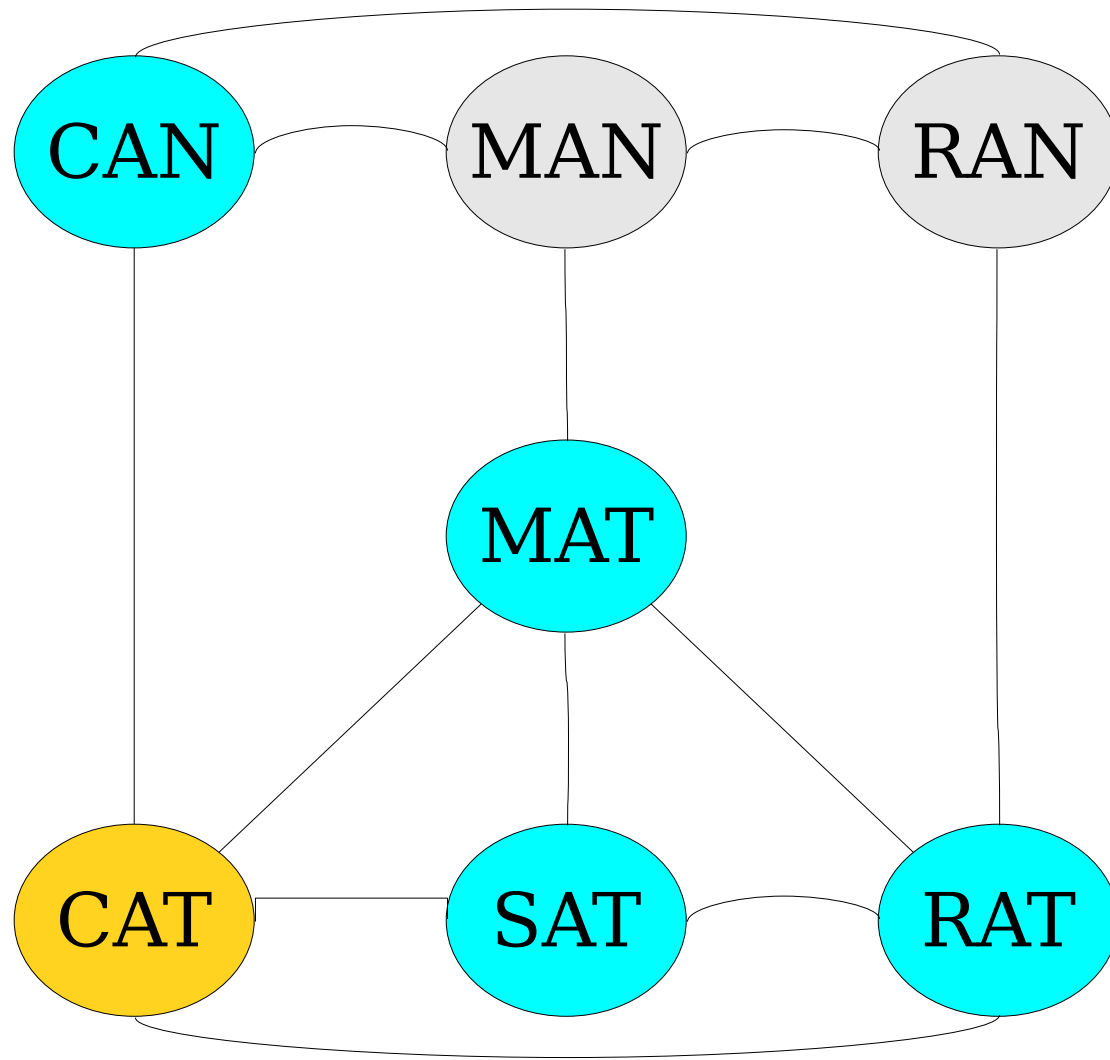
Two nodes are called *adjacent* if there is an edge between them.



nodes are called ***adjacent*** if there is an edge between them.



nodes are called ***adjacent*** if there is an edge between t



nodes are called ***adjacent*** if there is an edge between t

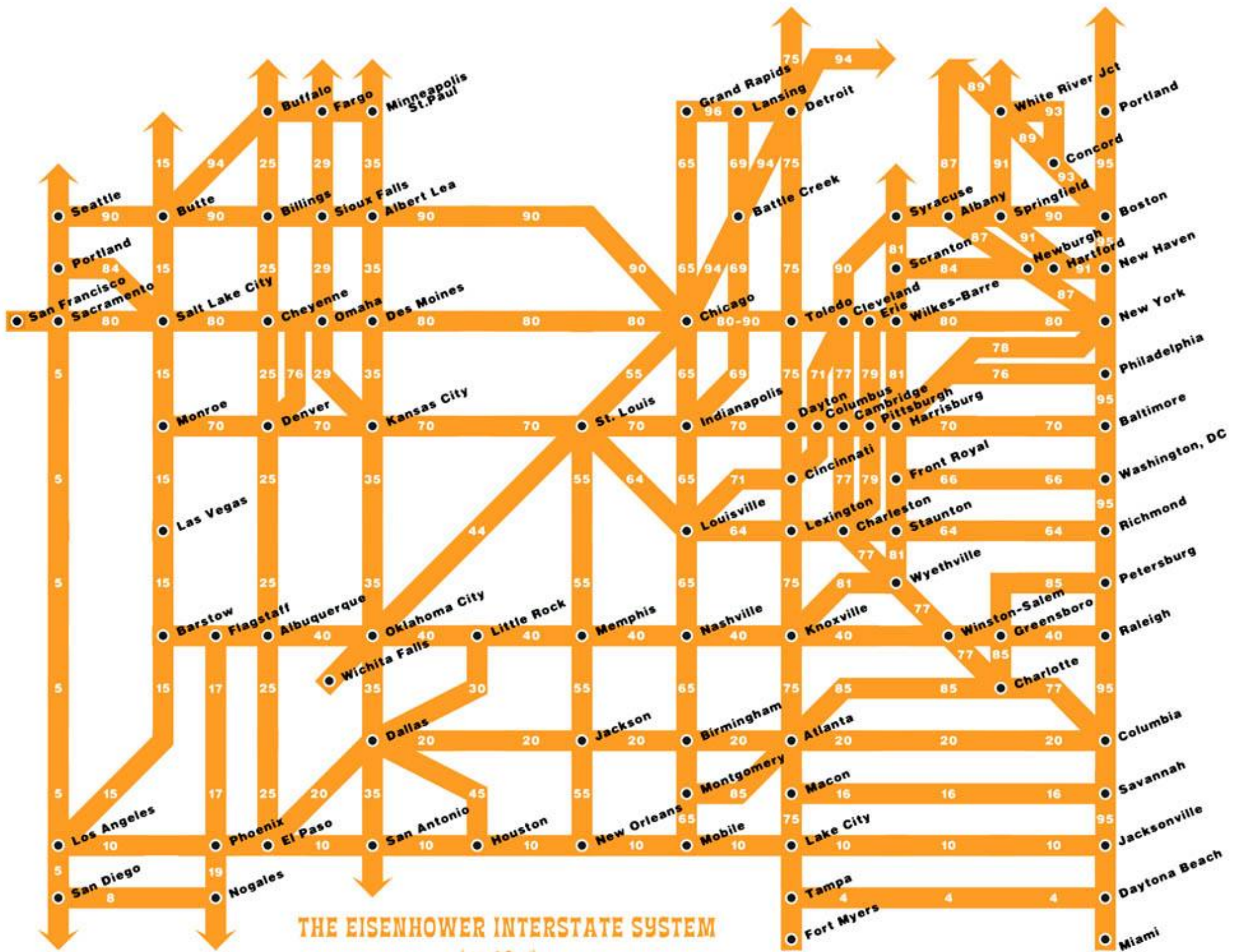


# Using our Formalisms

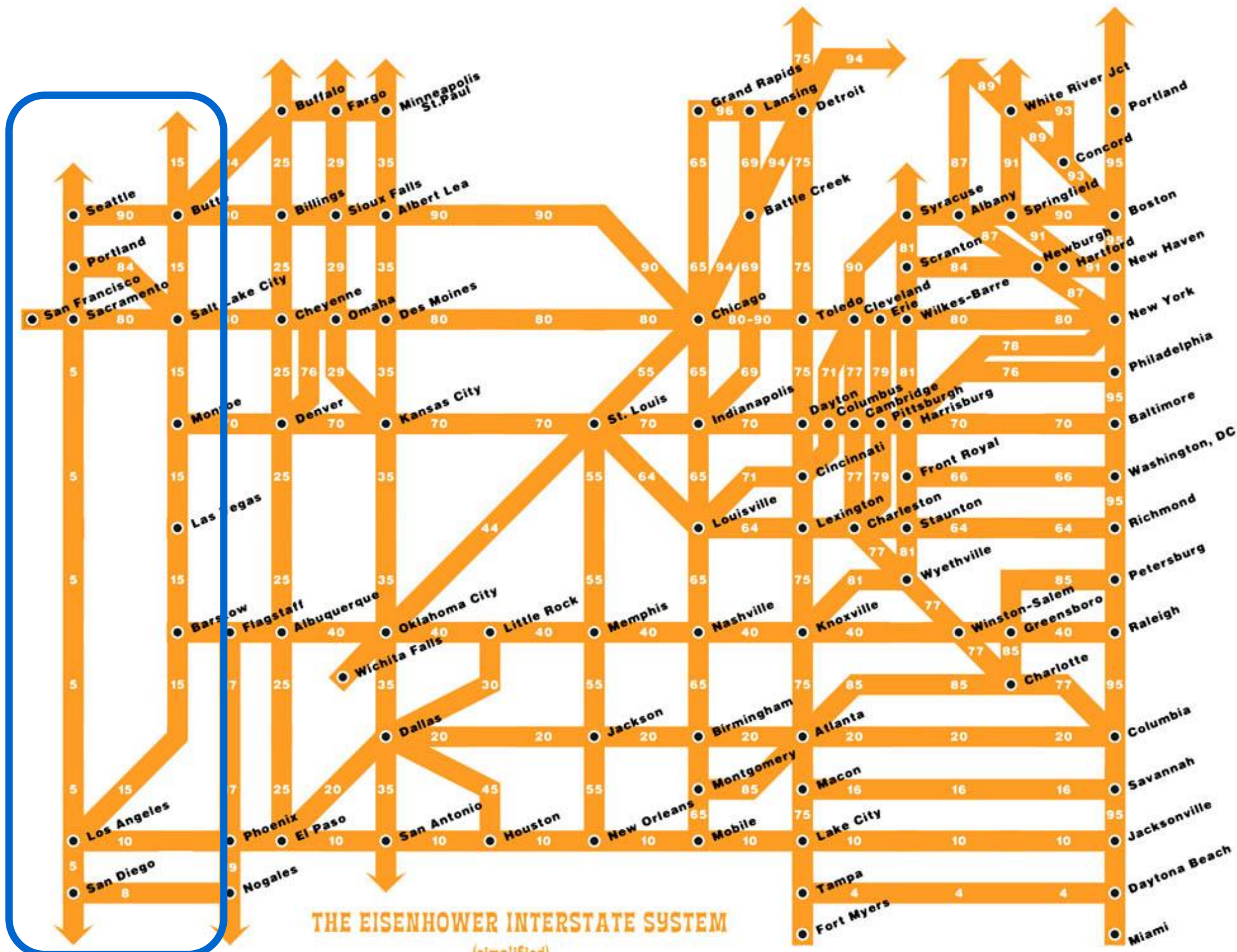
Let  $G = (V, E)$  be a graph.

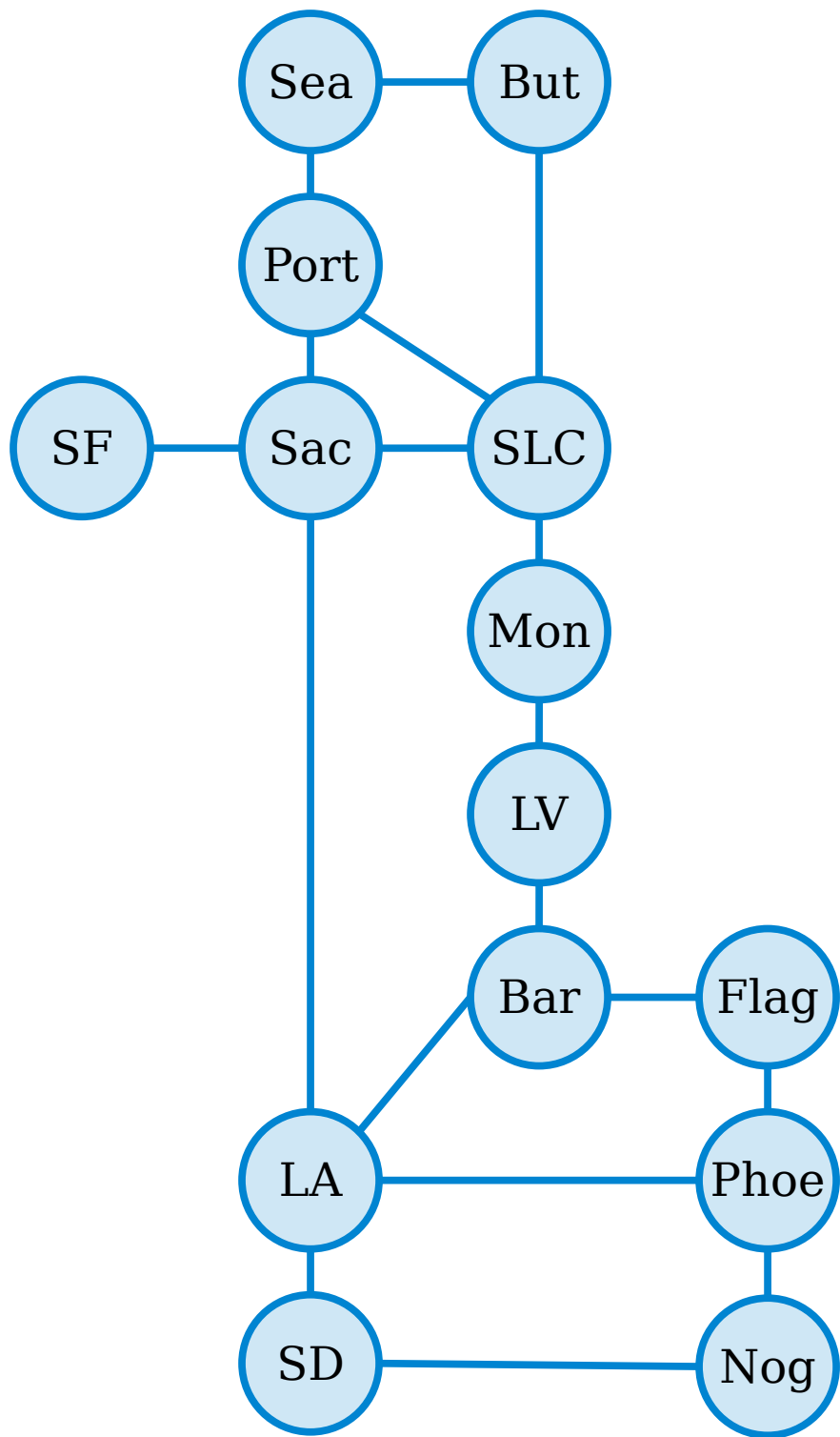
Intuitively, two nodes are adjacent if they're linked by an edge.

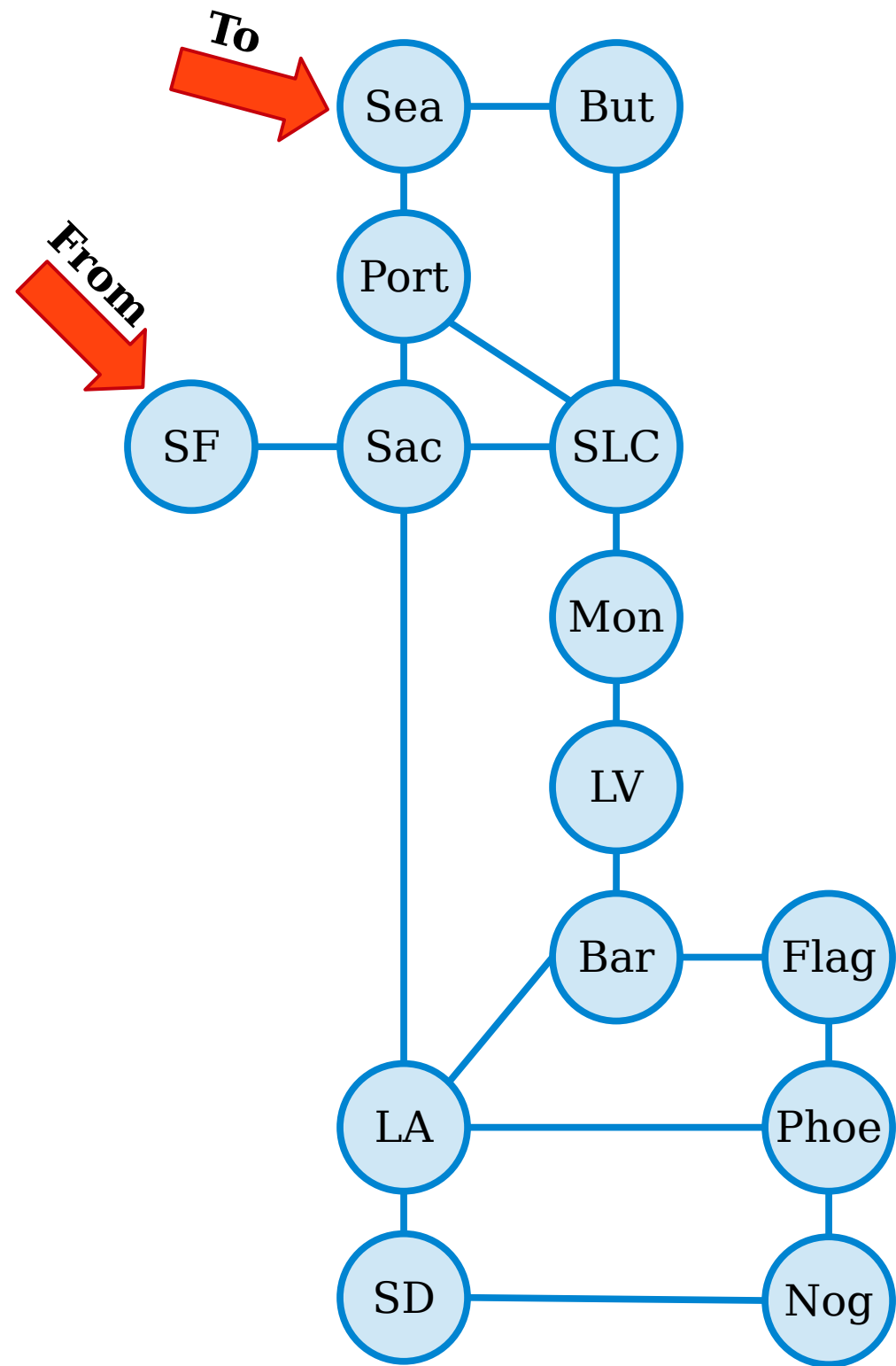
Formally speaking, we say that two nodes  $u, v \in V$  are **adjacent** if  $\{u, v\} \in E$ .

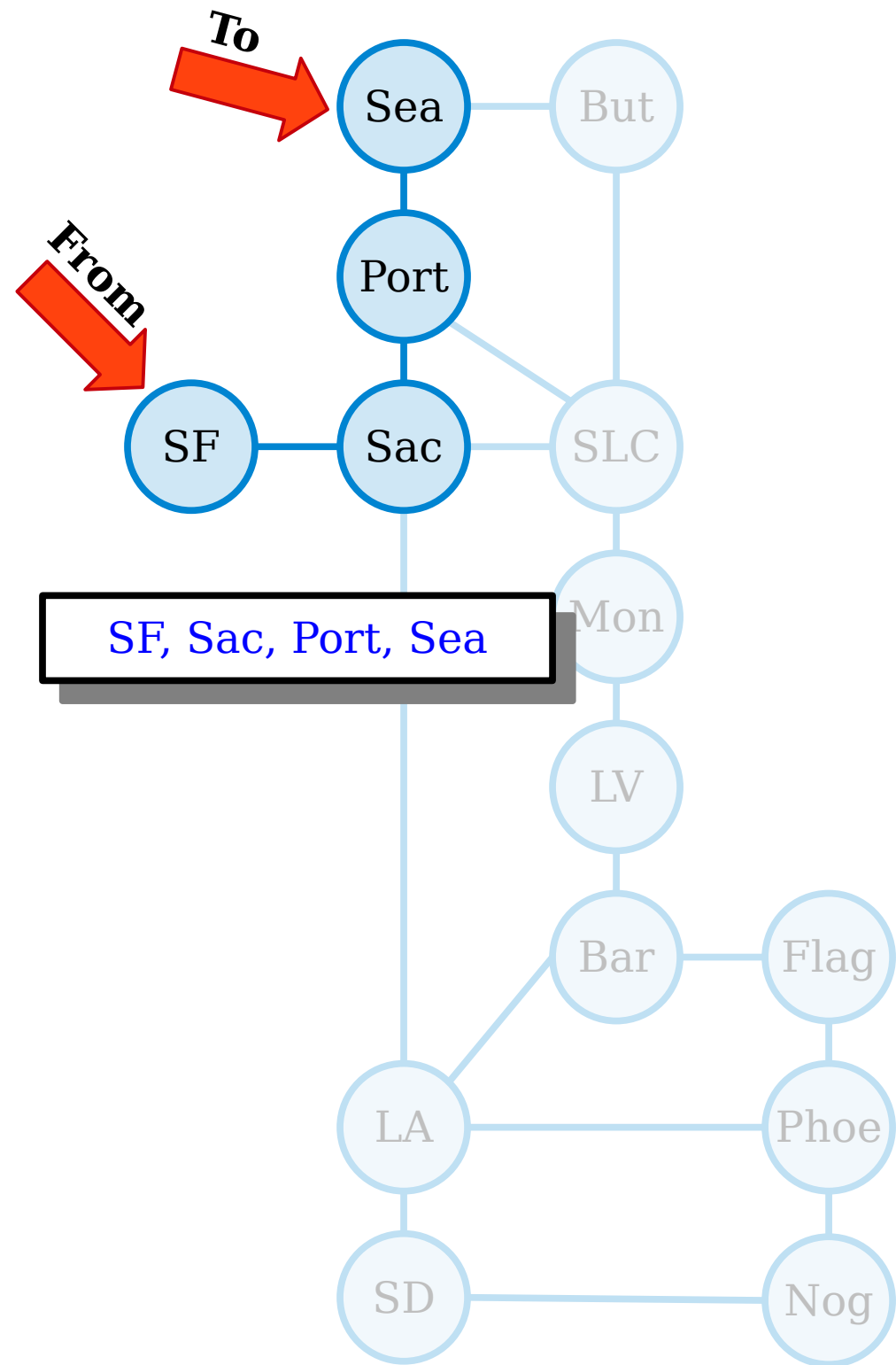




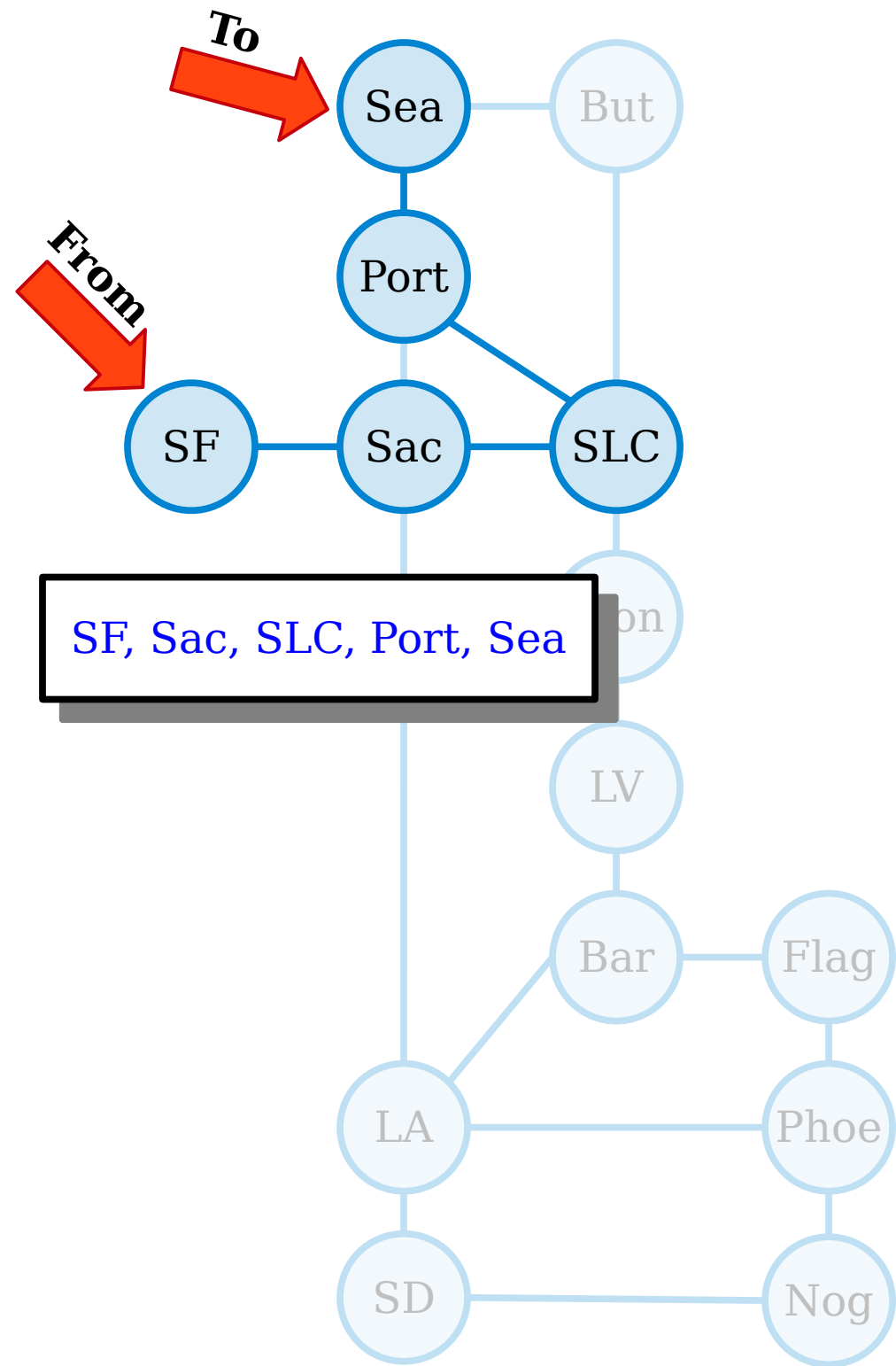


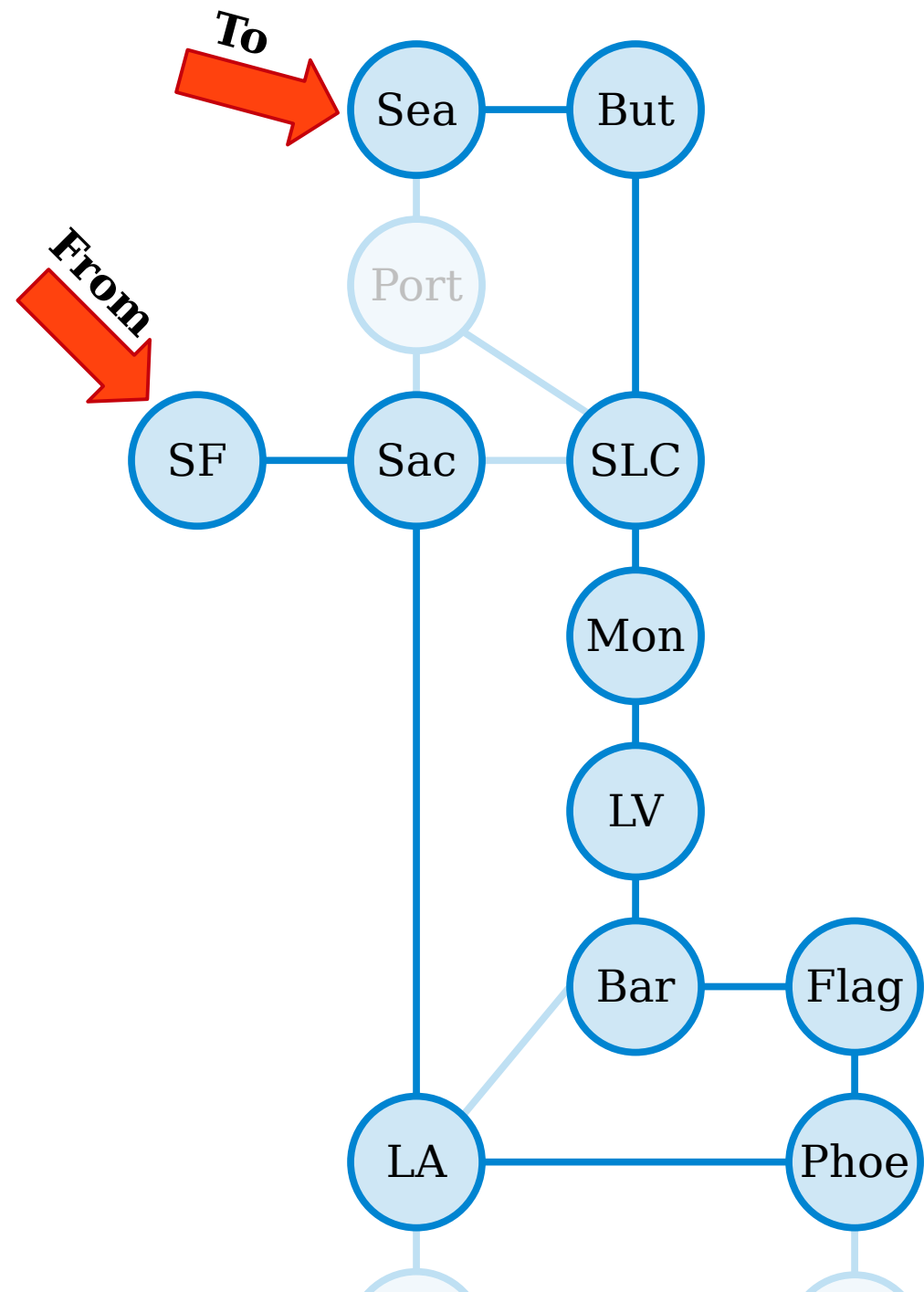






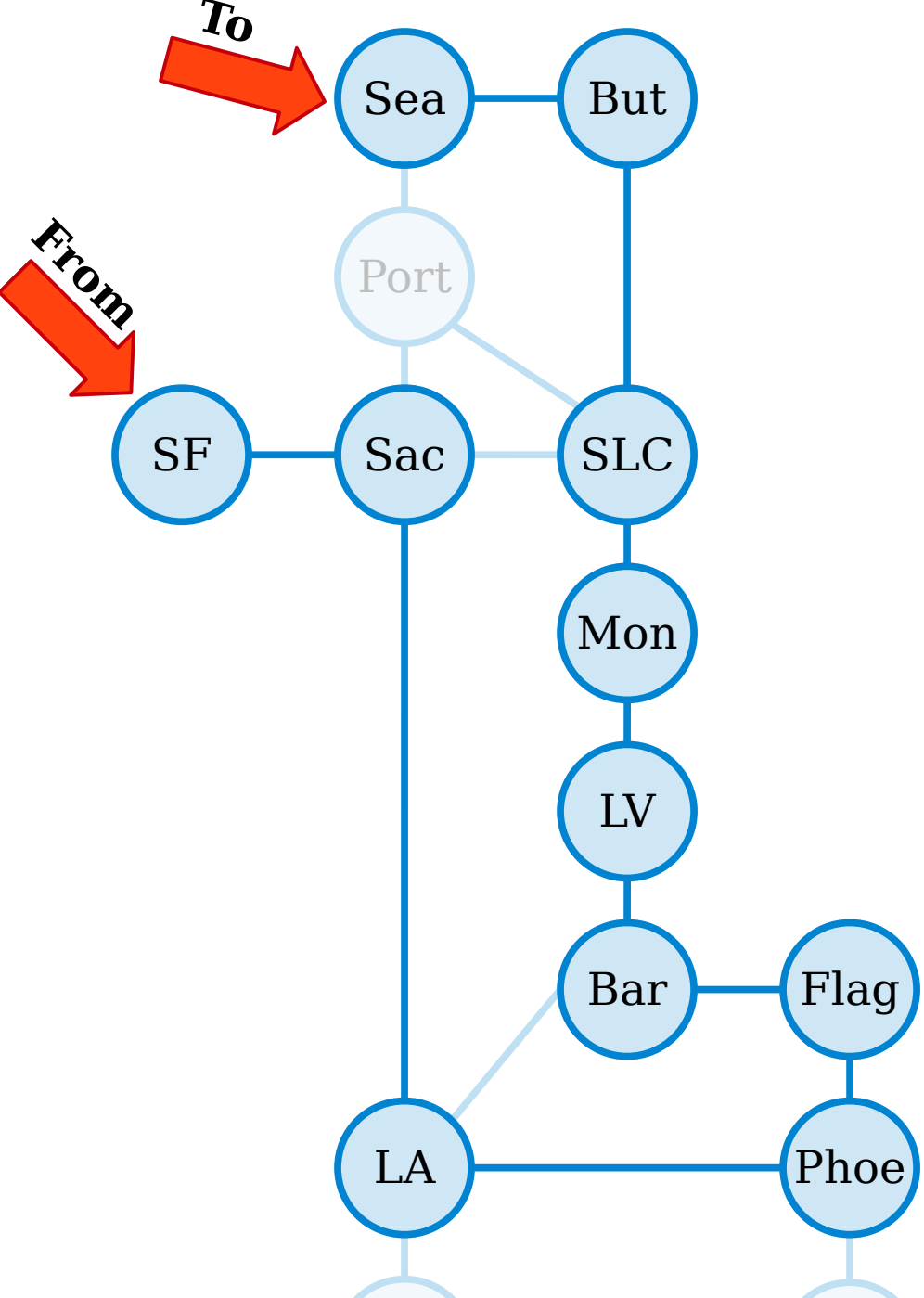




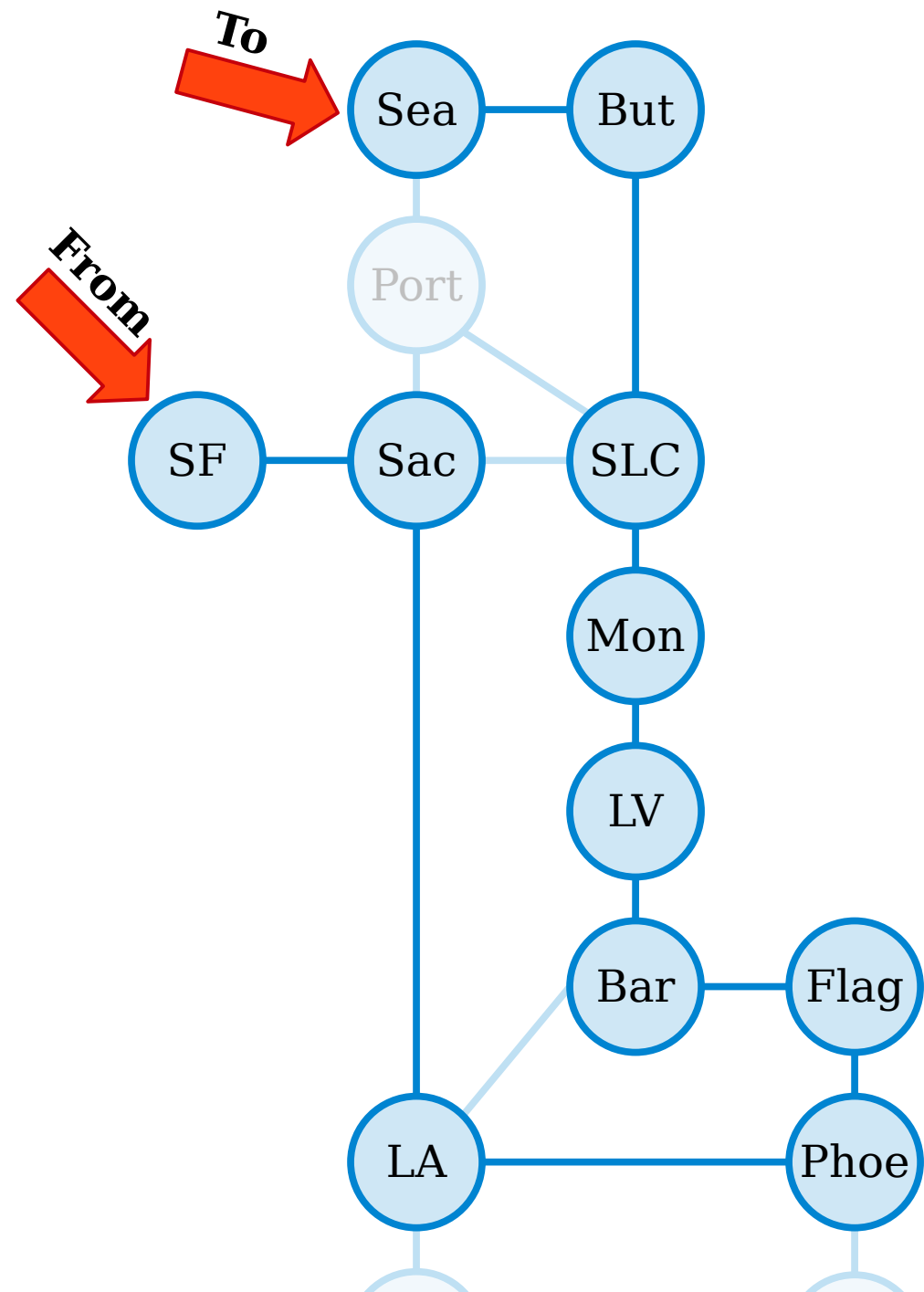


SF, Sac, LA, Phoe, Flag, Bar, LV, Mon, SLC, But, Sea

A **path** in a graph  $G = (V, E)$  is a sequence of one or more nodes  $v_1, v_2, v_3, \dots, v_n$  such that any two consecutive nodes in the sequence are adjacent.



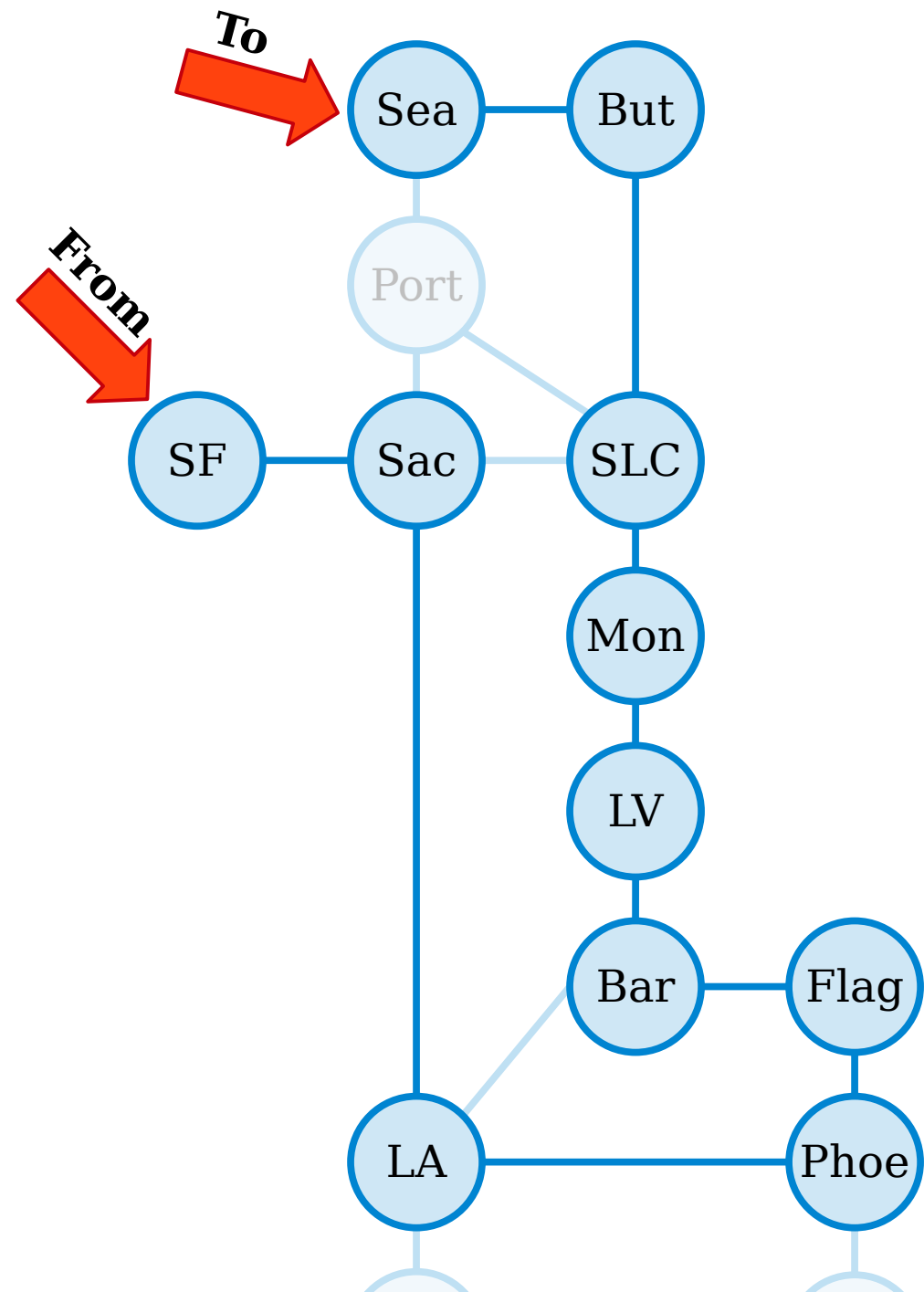
SF, Sac, LA, Phoe, Flag, Bar, LV, Mon, SLC, But, Sea



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The **length** of the path  $v_1, \dots, v_n$  is  $n - 1$ .

SF, Sac, LA, Phoe, Flag, Bar, LV, Mon, SLC, But, Sea

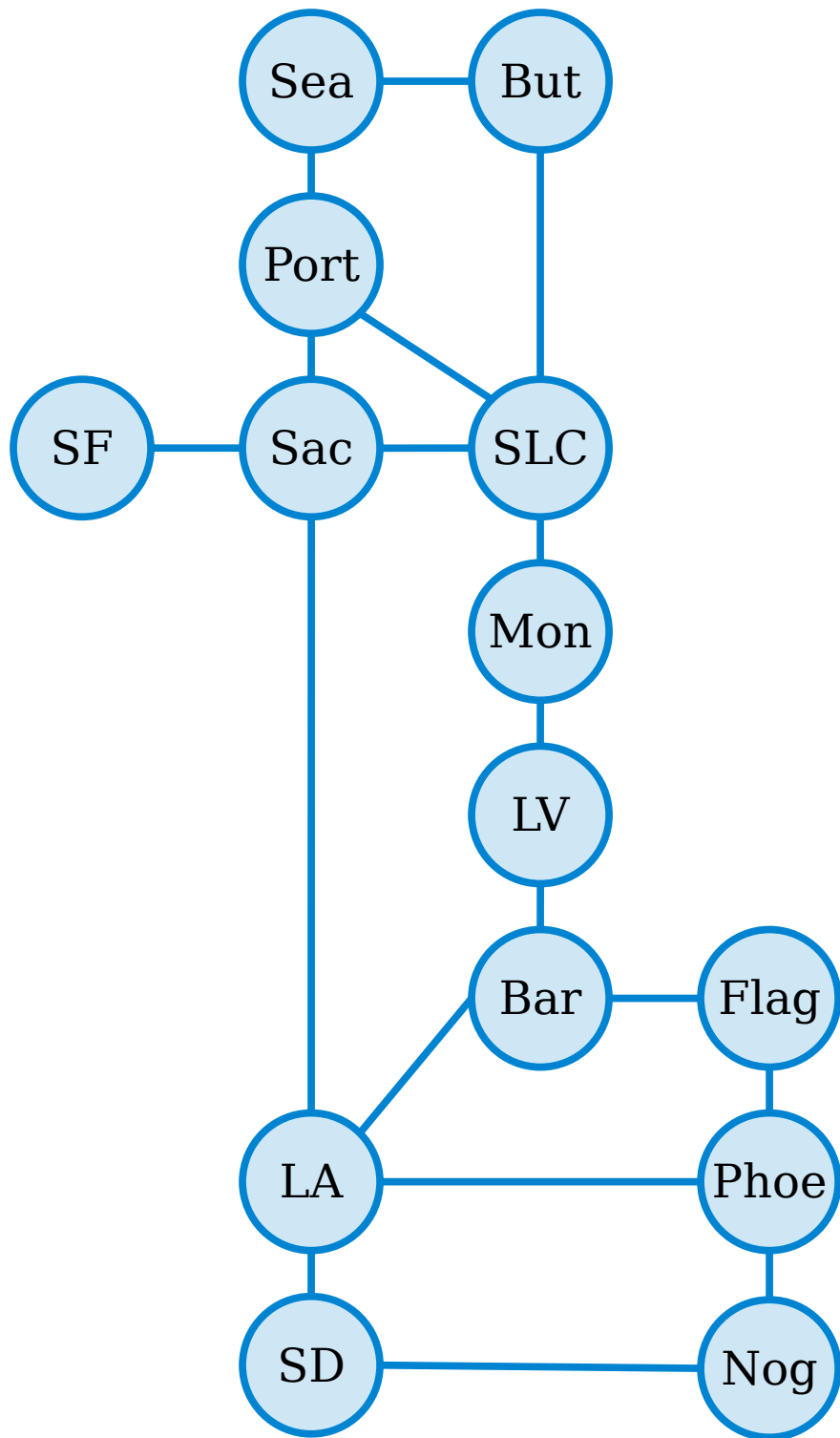


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(This path has length 10, but visits 11 cities.)

SF, Sac, LA, Phoe, Flag, Bar, LV, Mon, SLC, But, Sea

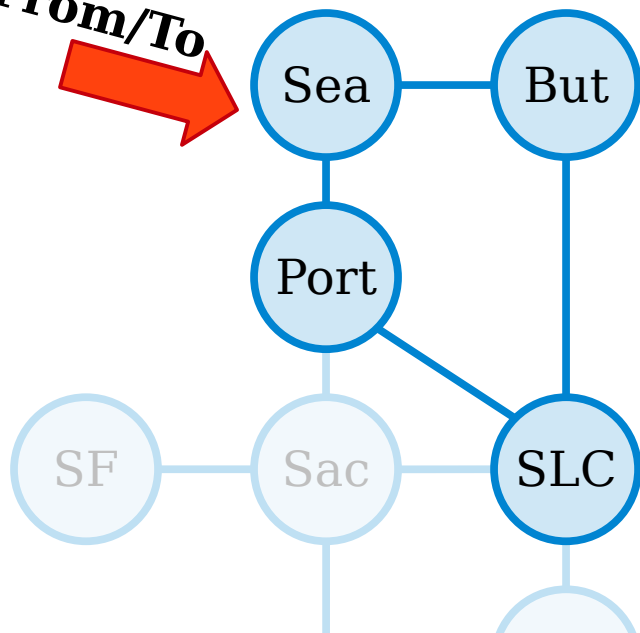


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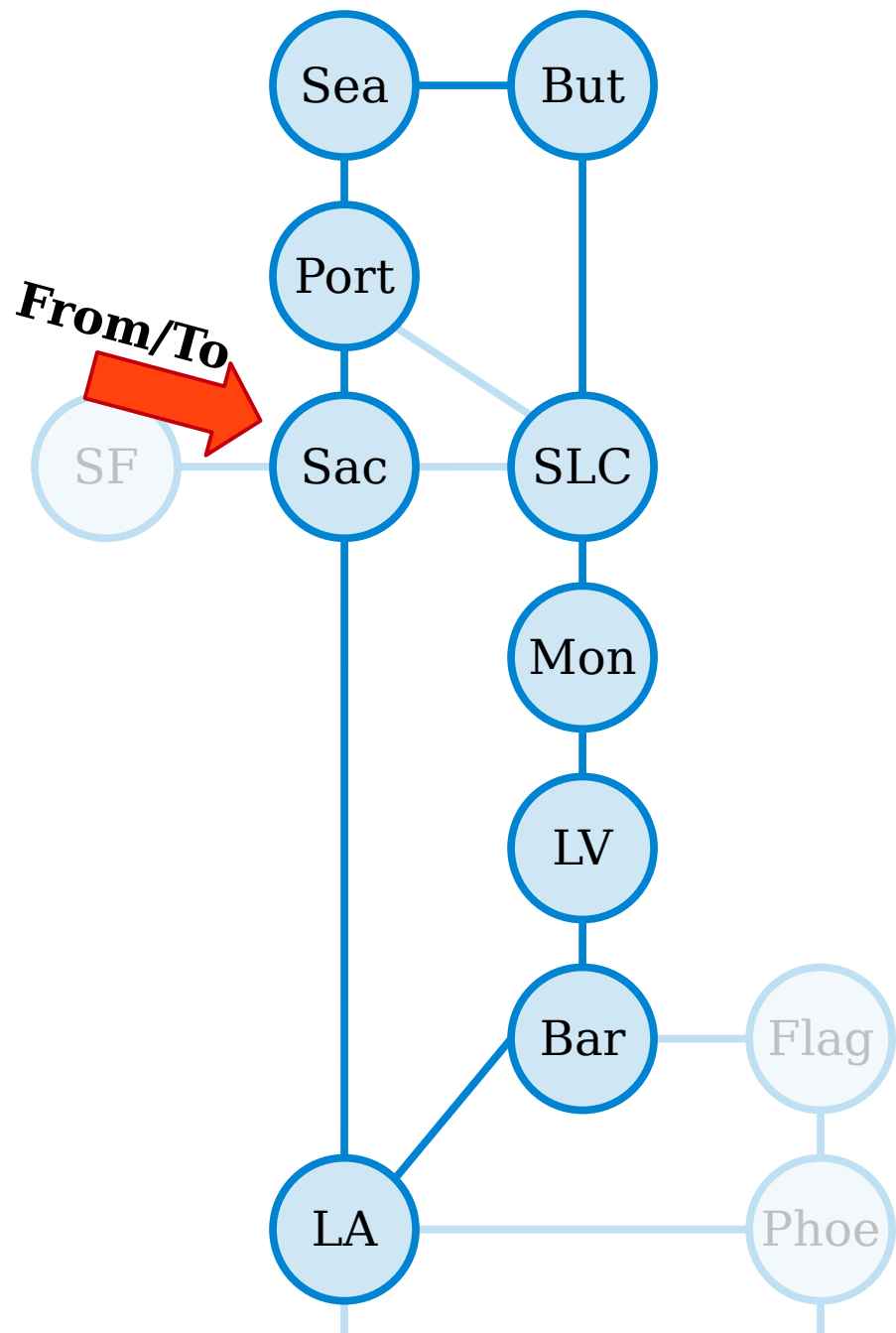
**From/To**



Sea, But, SLC, Port, Sea

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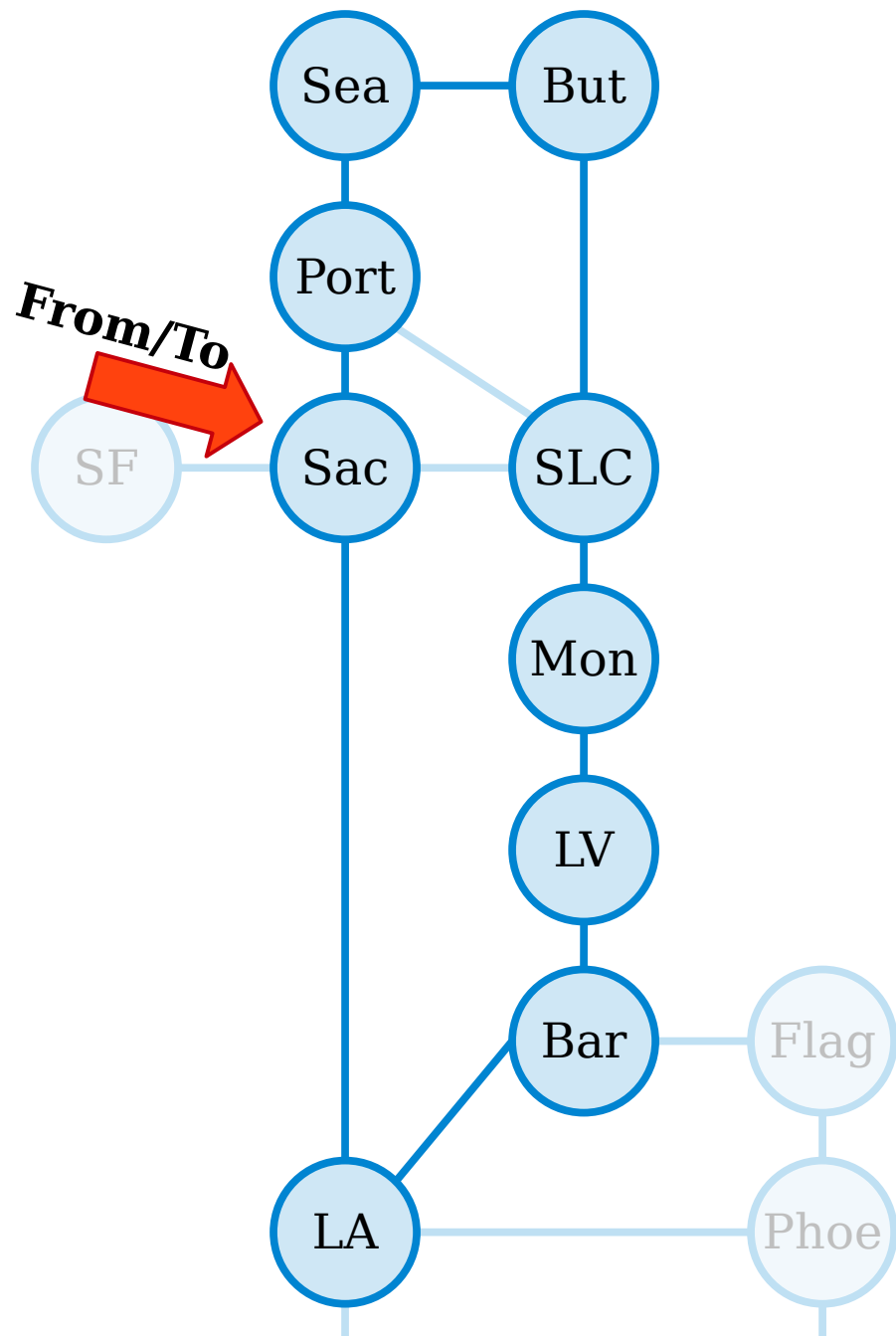
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Sac, Port, Sea, But, SLC, Mon, LV, Bar, LA, Sac

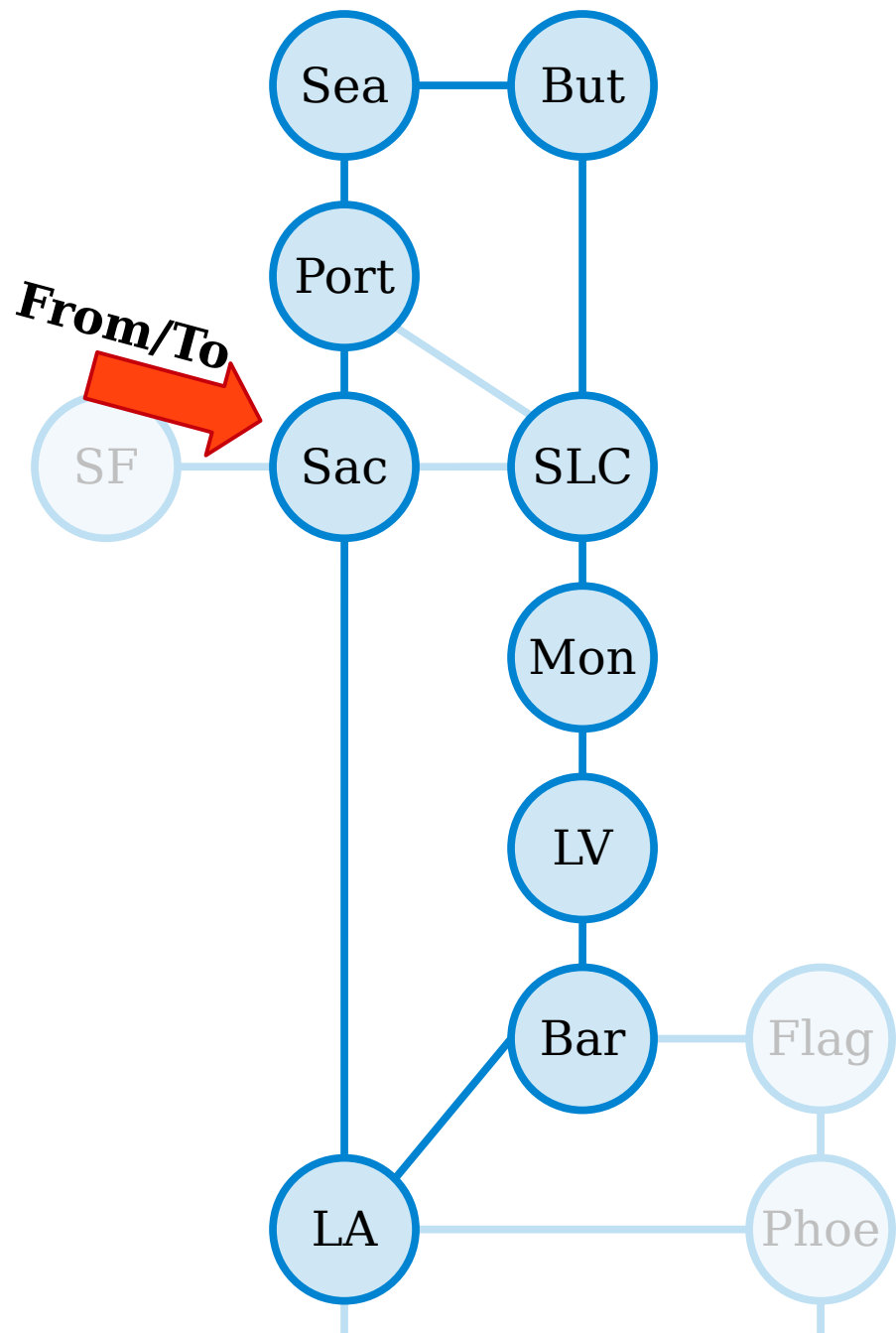


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A **cycle** in a graph is a path from a node back to itself. (By convention, a cycle cannot have length zero.)

Sac, Port, Sea, But, SLC, Mon, LV, Bar, LA, Sac



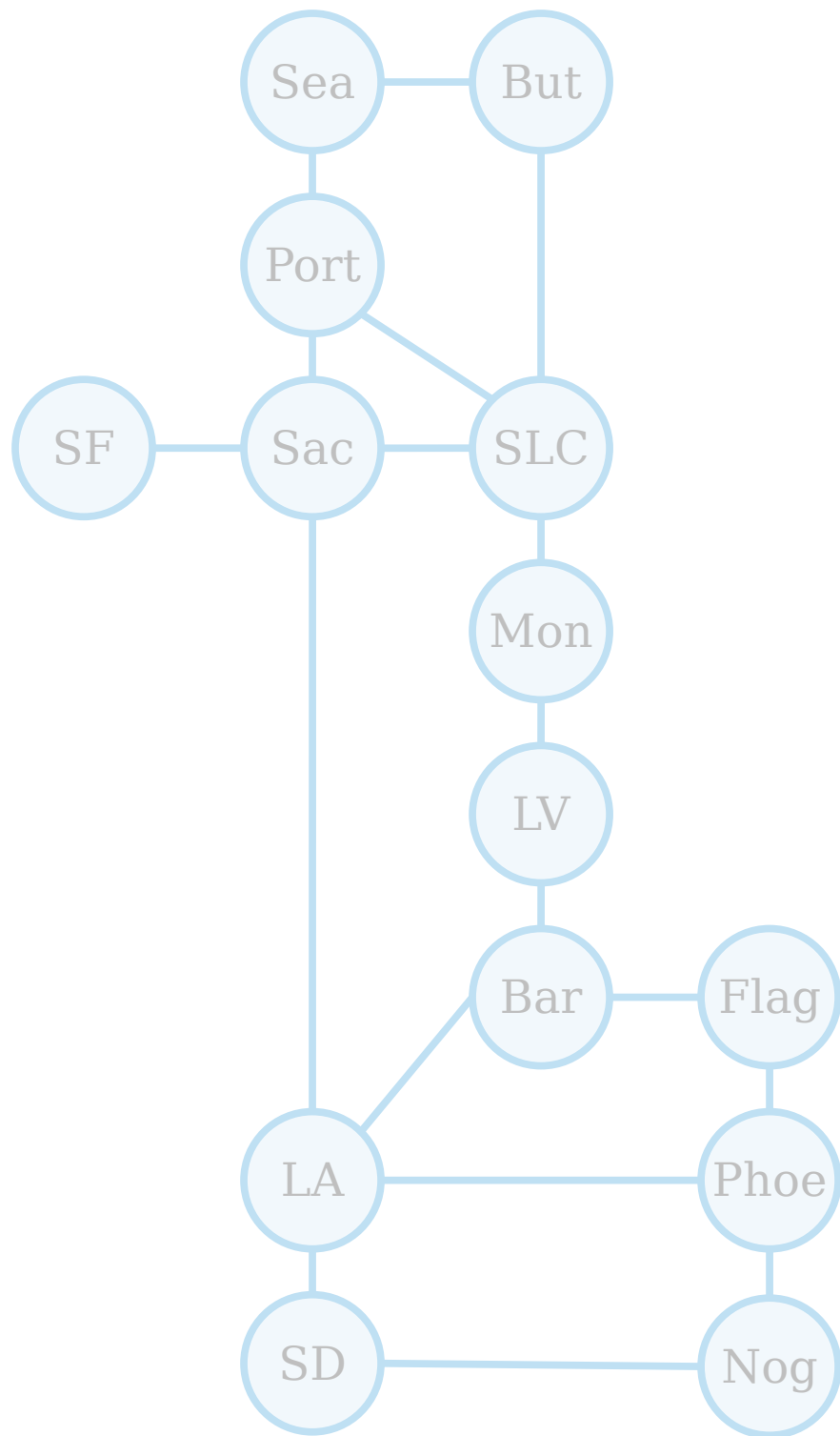
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(This cycle has length nine and visits nine different cities.)

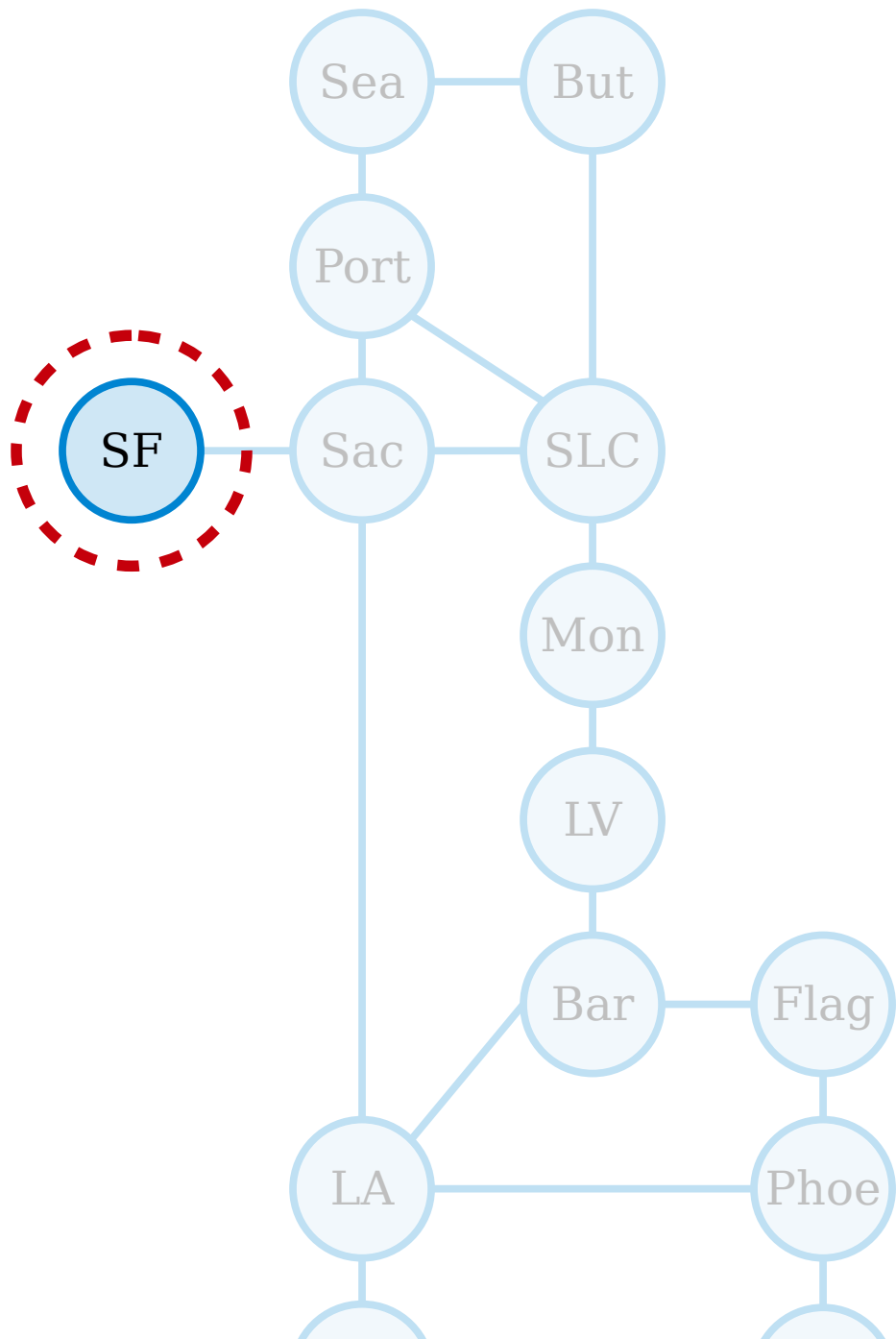
Sac, Port, Sea, But, SLC, Mon, LV, Bar, LA, Sac



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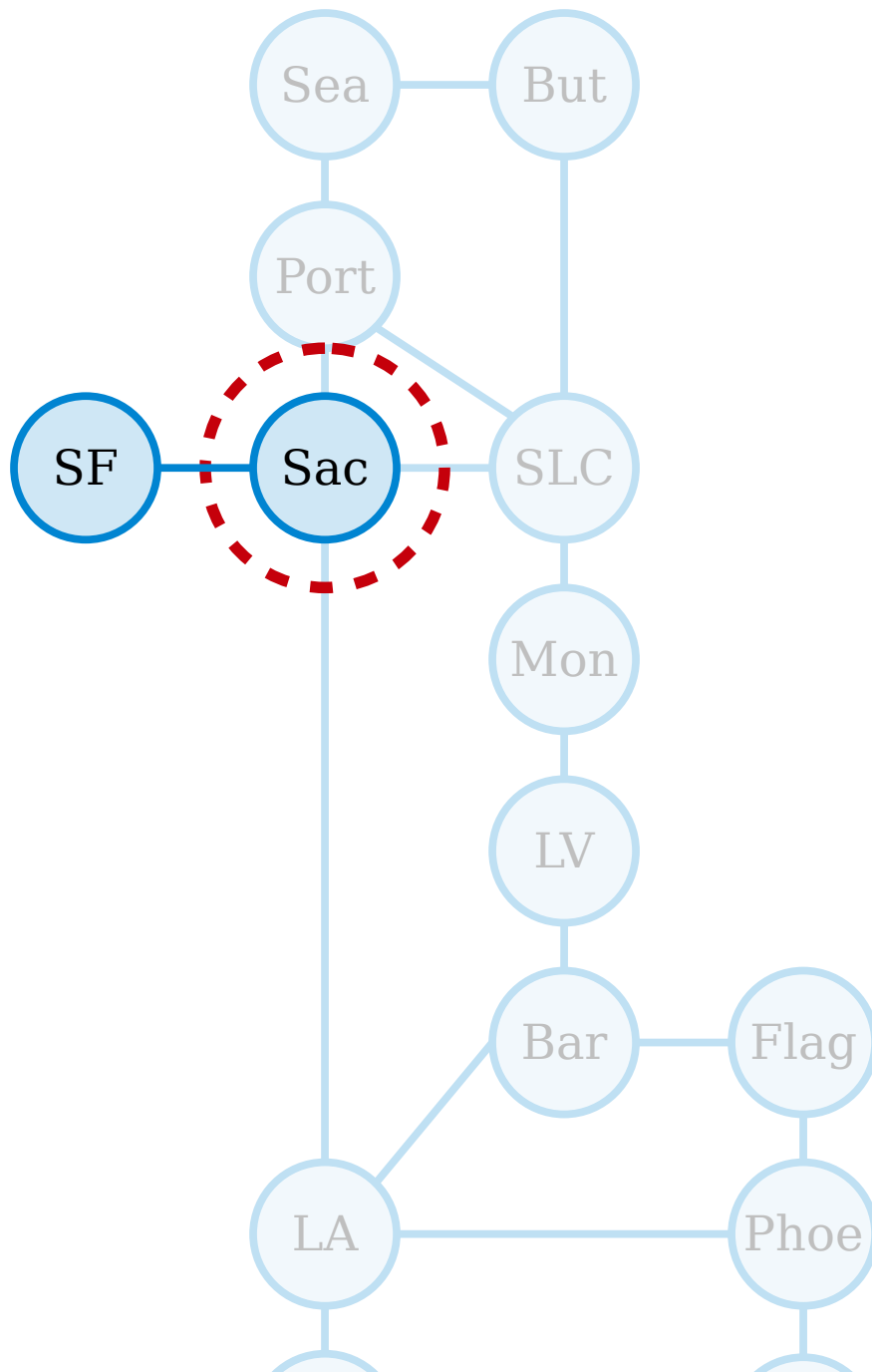
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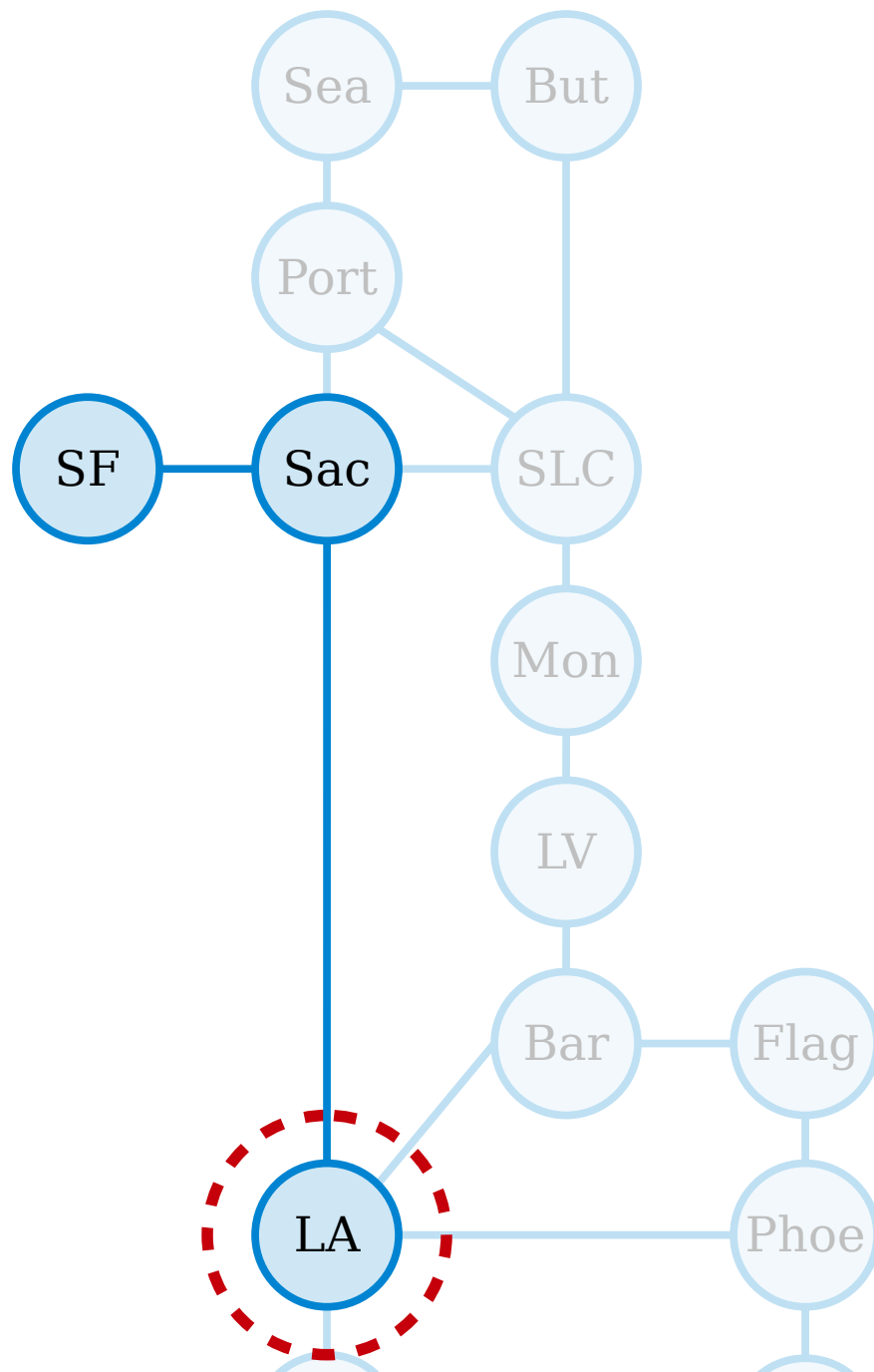
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SF, Sac



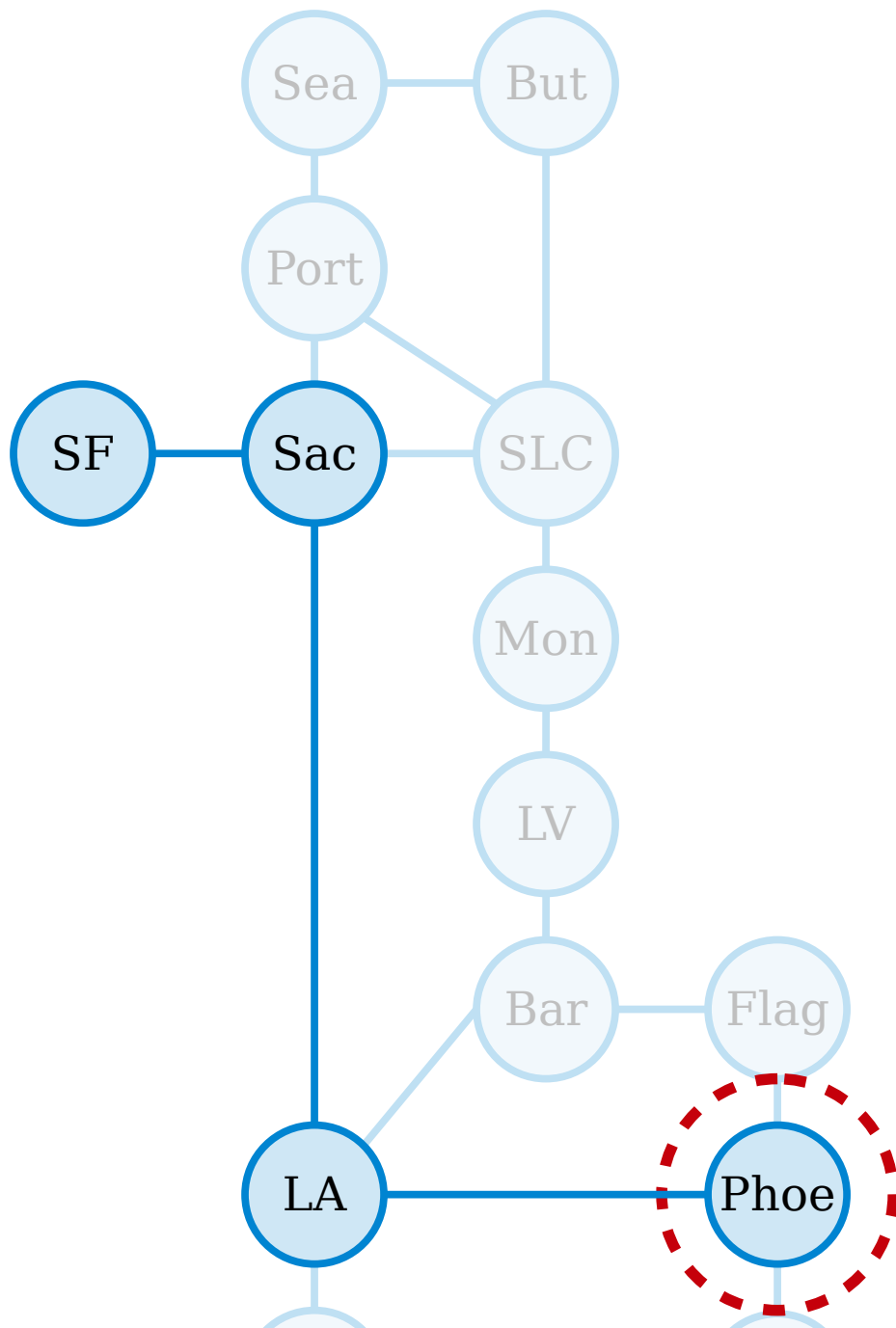


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SF, Sac, LA

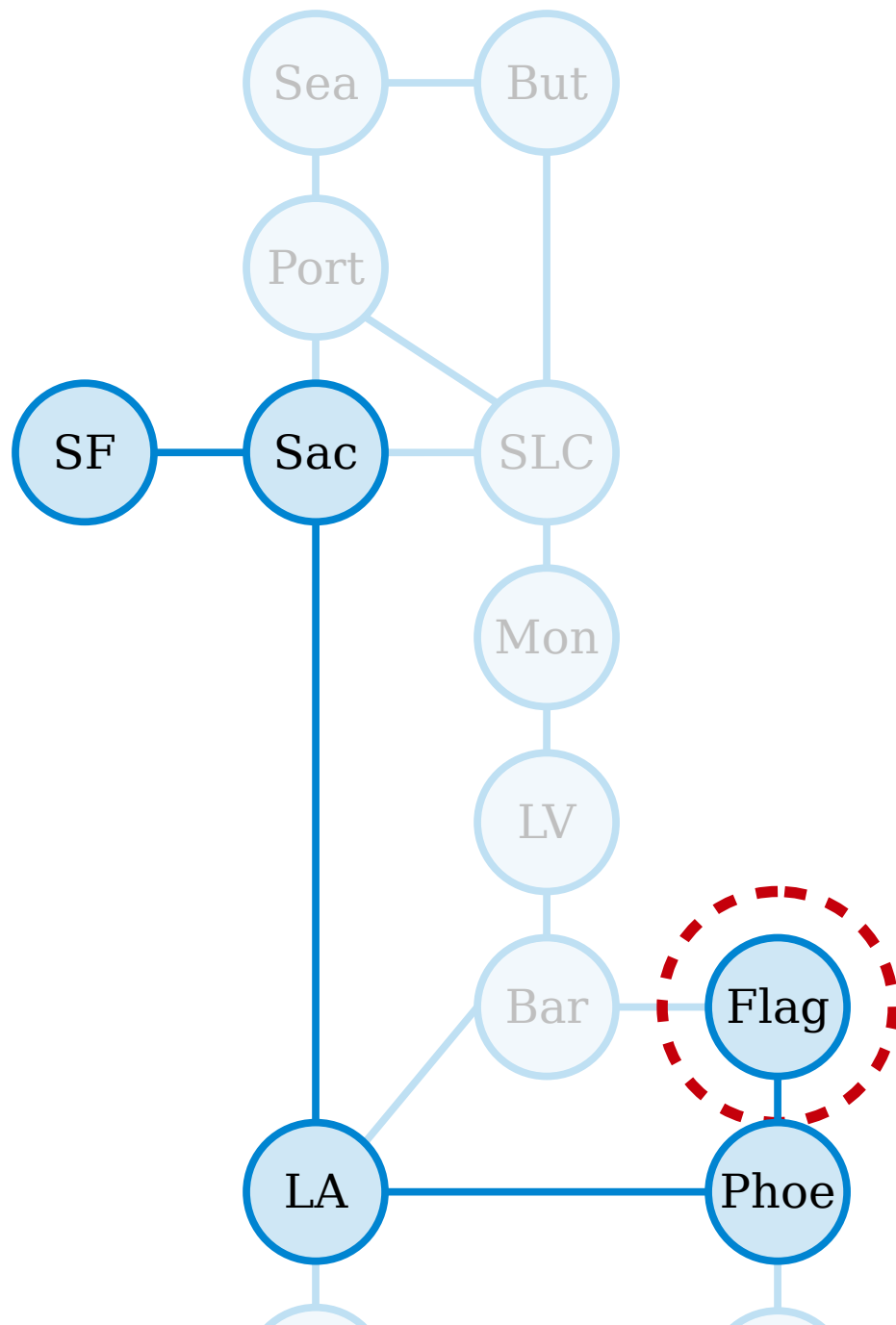


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SF, Sac, LA, Phoe

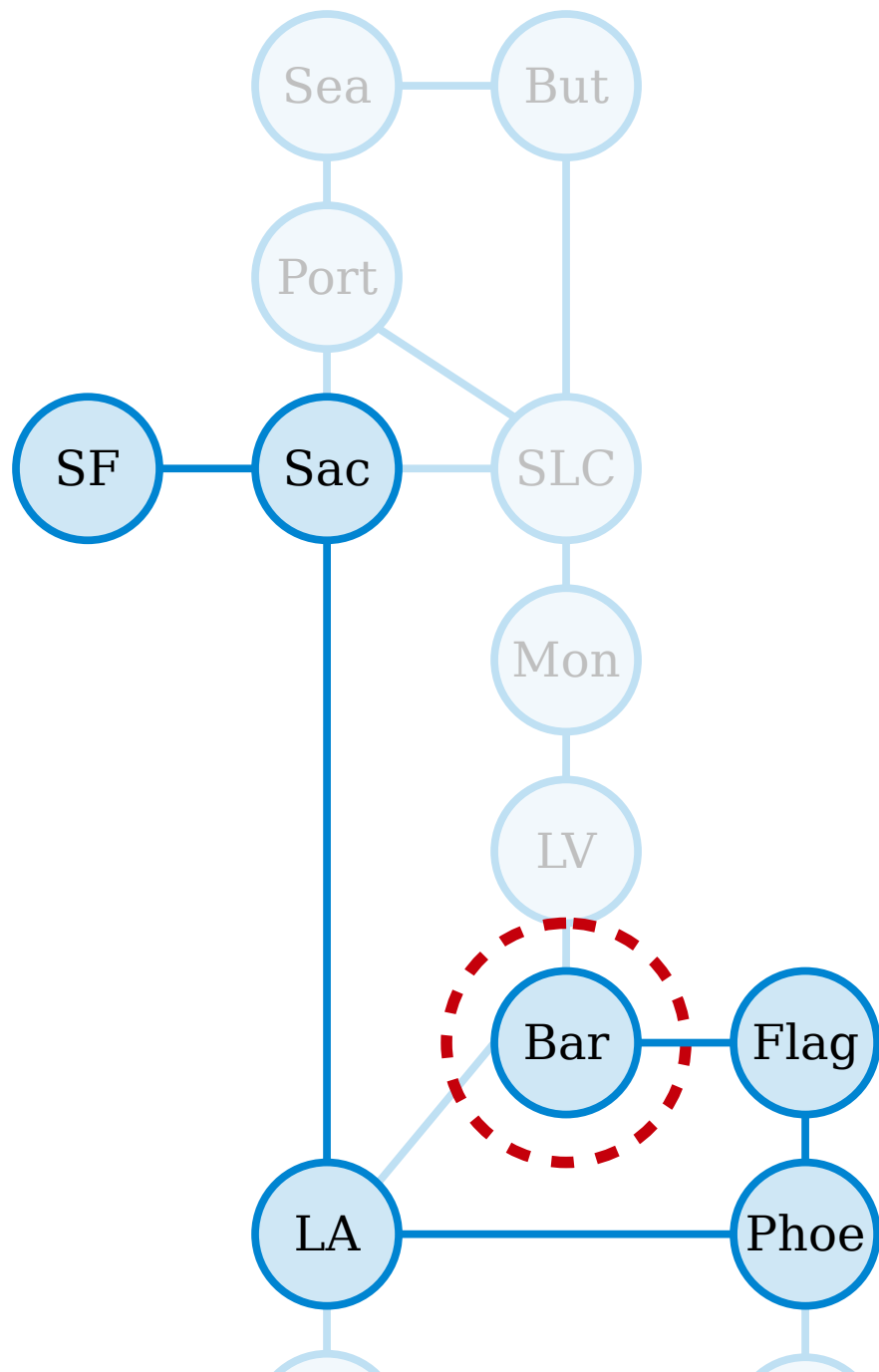


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SF, Sac, LA, Phoe, Flag

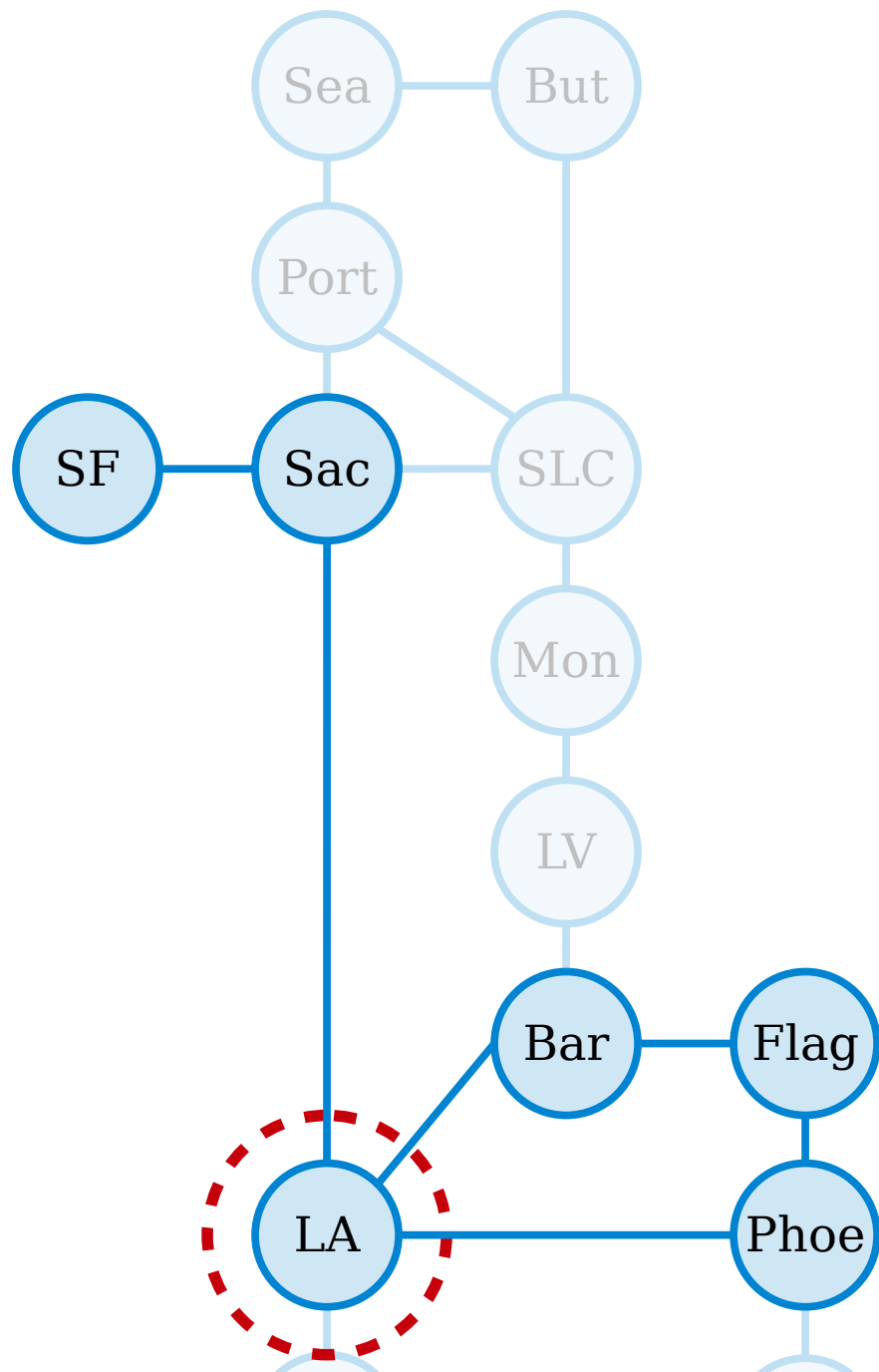


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SF, Sac, LA, Phoe, Flag, Bar

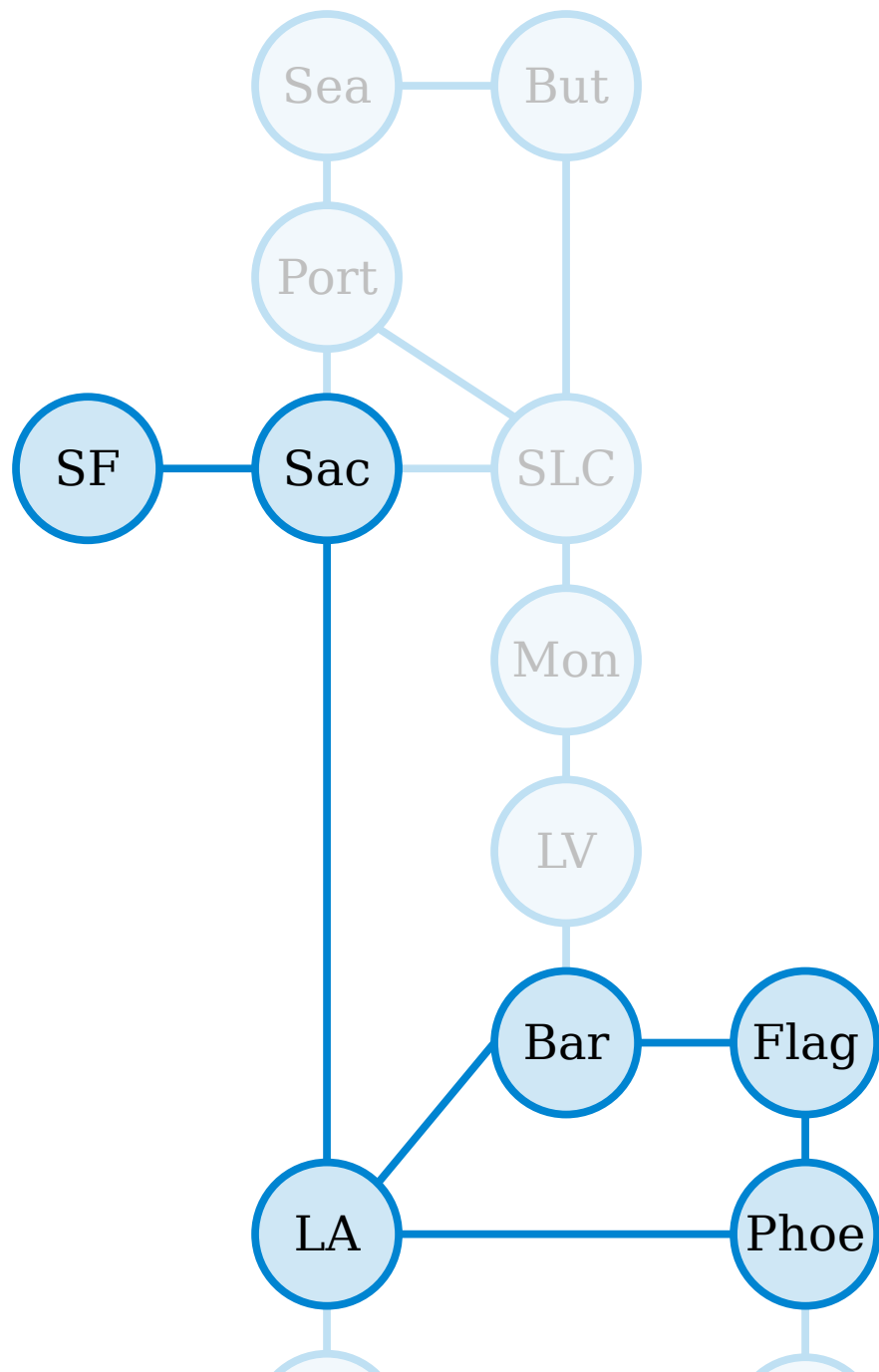


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SF, Sac, LA, Phoe, Flag, Bar, LA

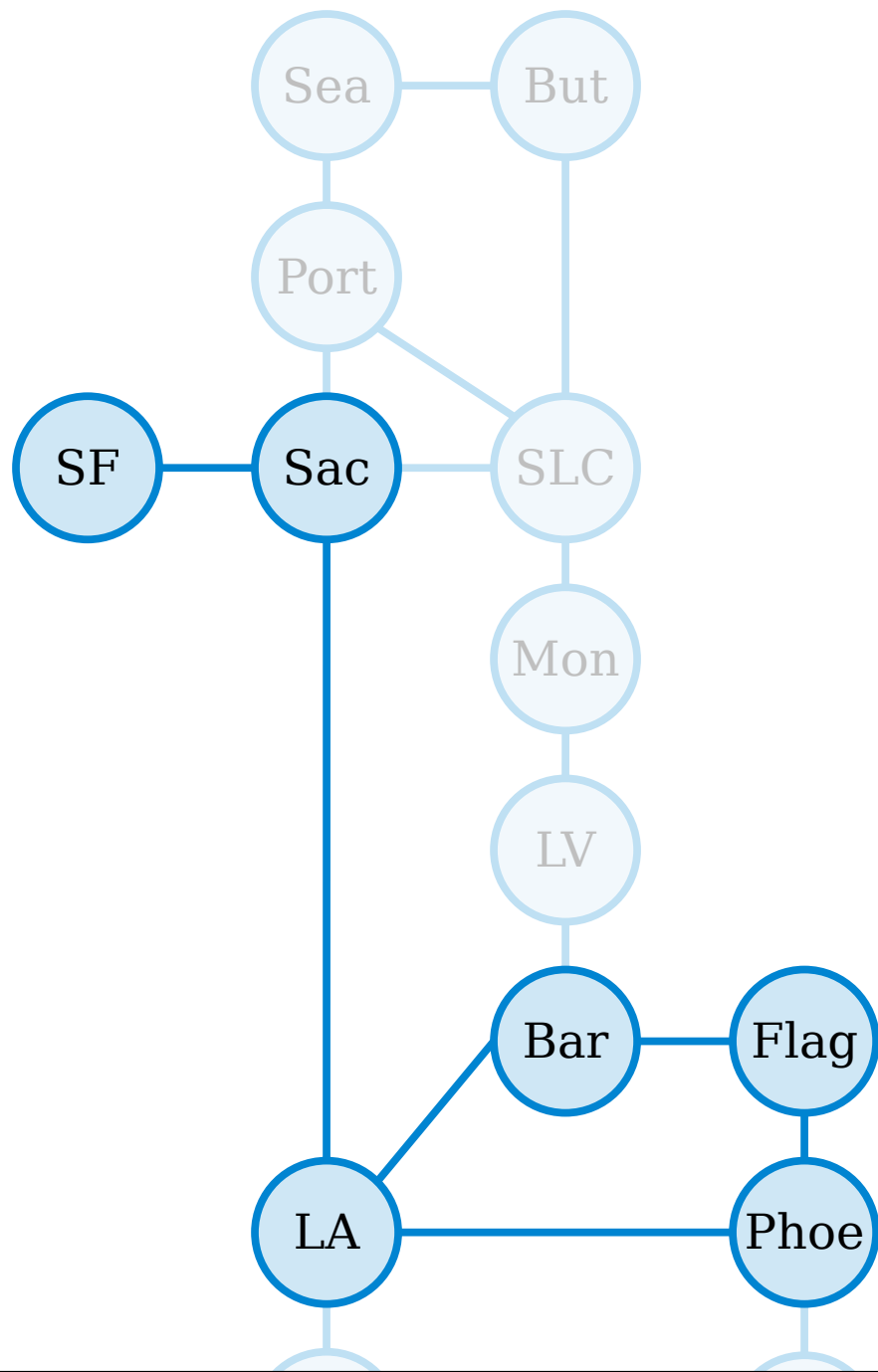


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SF, Sac, LA, Phoe, Flag, Bar, LA



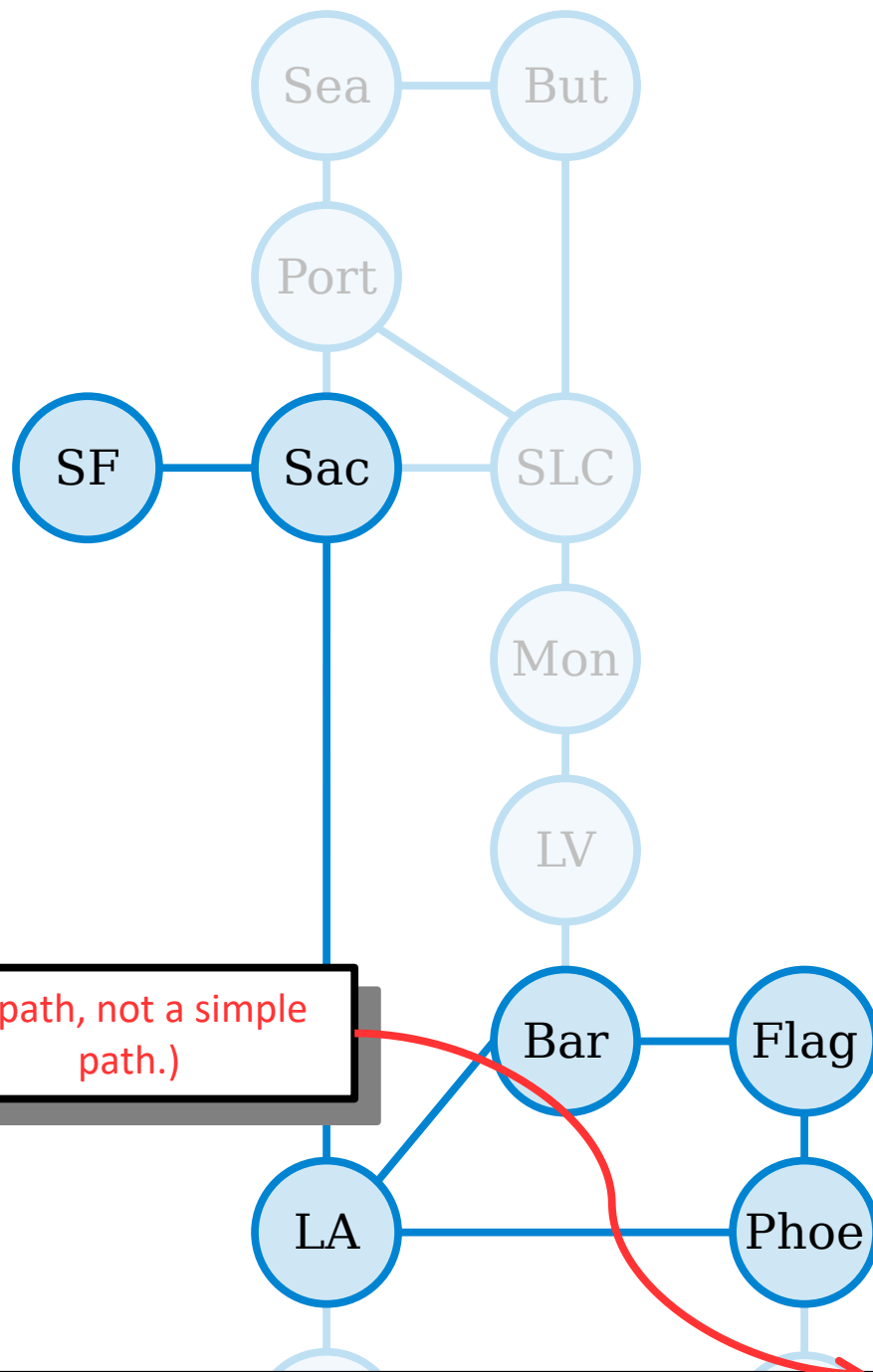
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A **simple path** in a graph is path that does not repeat any nodes or edges.

SF, Sac, LA, Phoe, Flag, Bar, LA



(A path, not a simple path.)

SF, Sac, LA, Phoe, Flag, Bar, LA

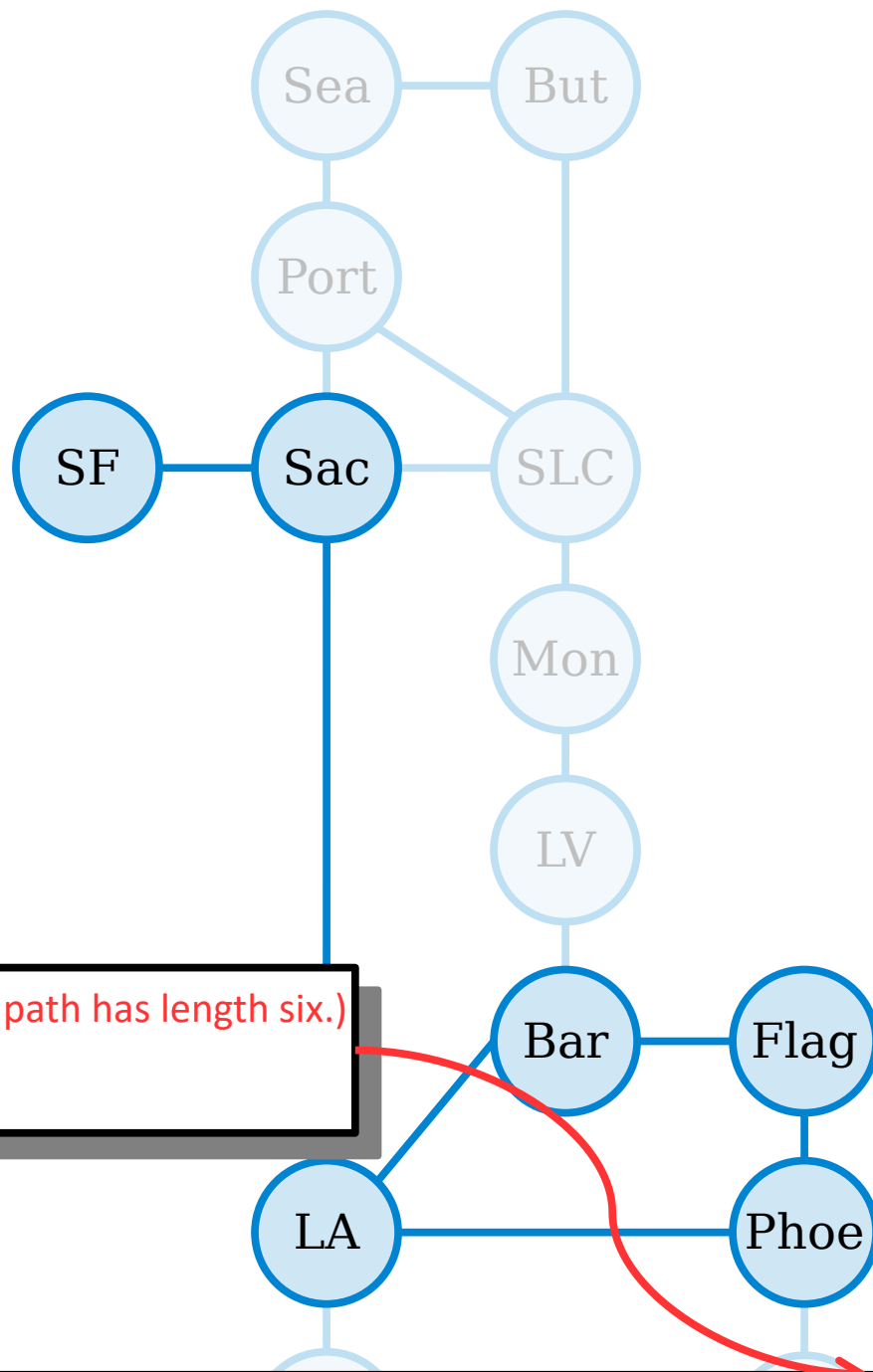
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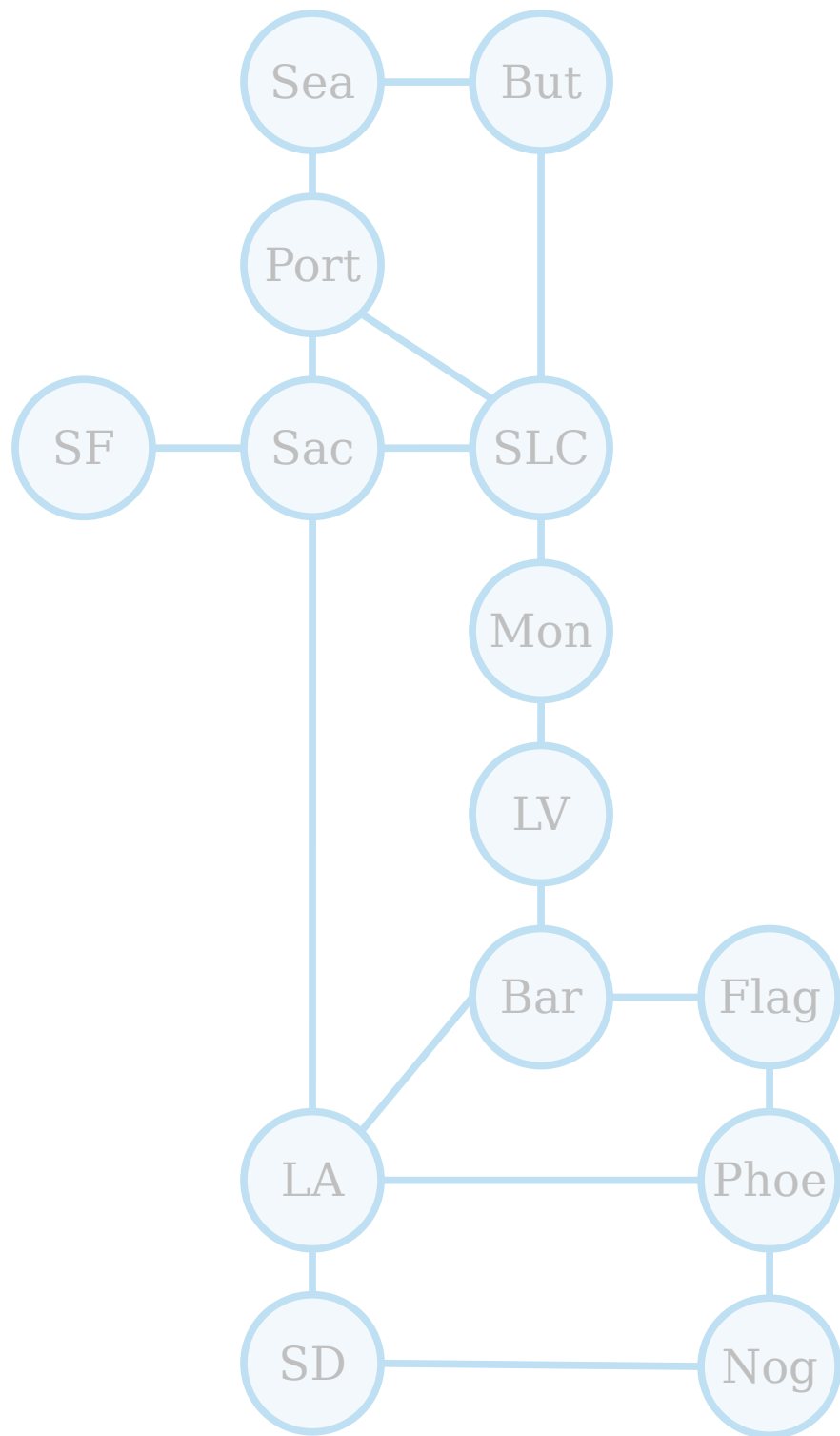
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SF, Sac, LA, Phoe, Flag, Bar, LA

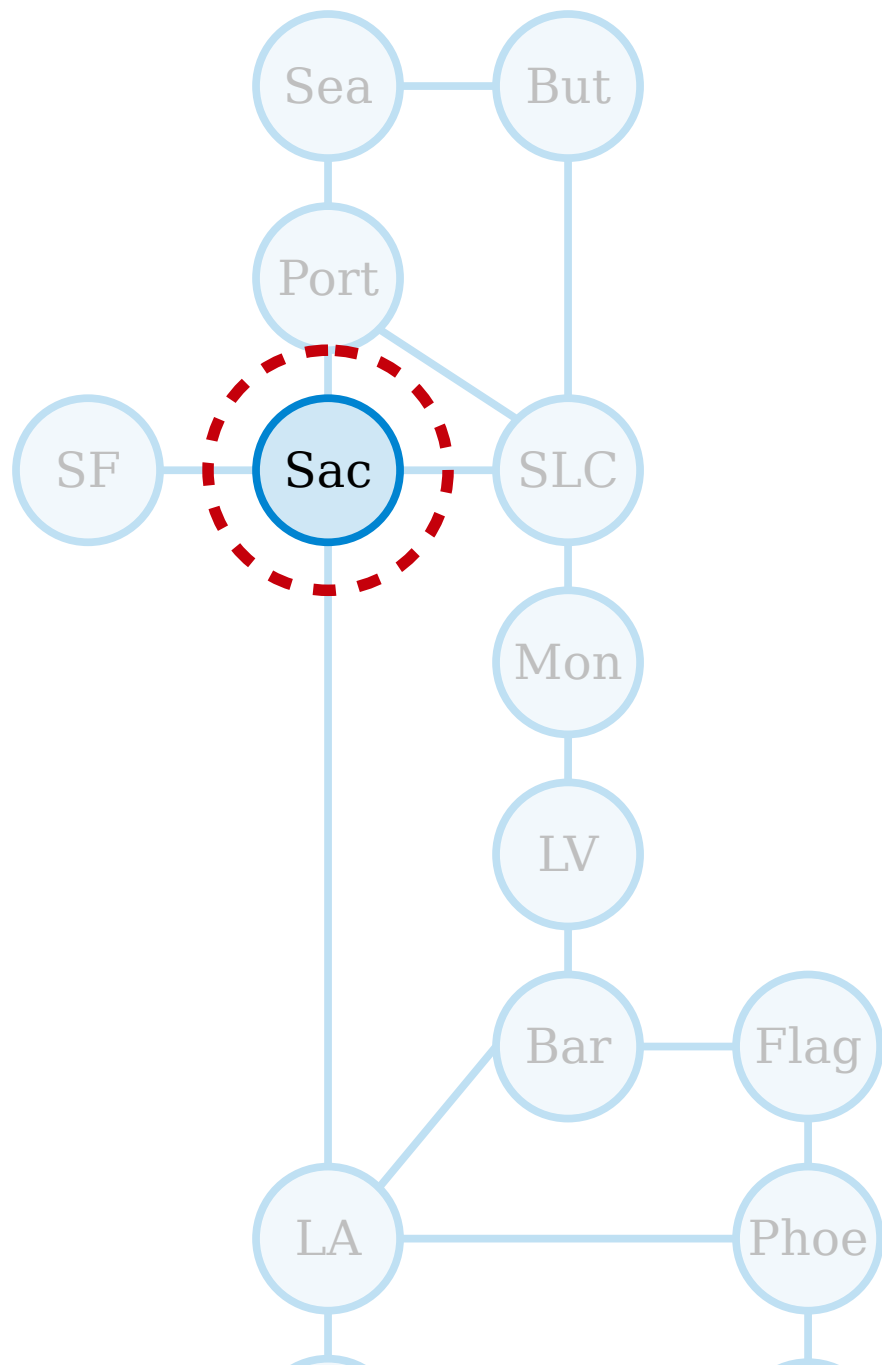


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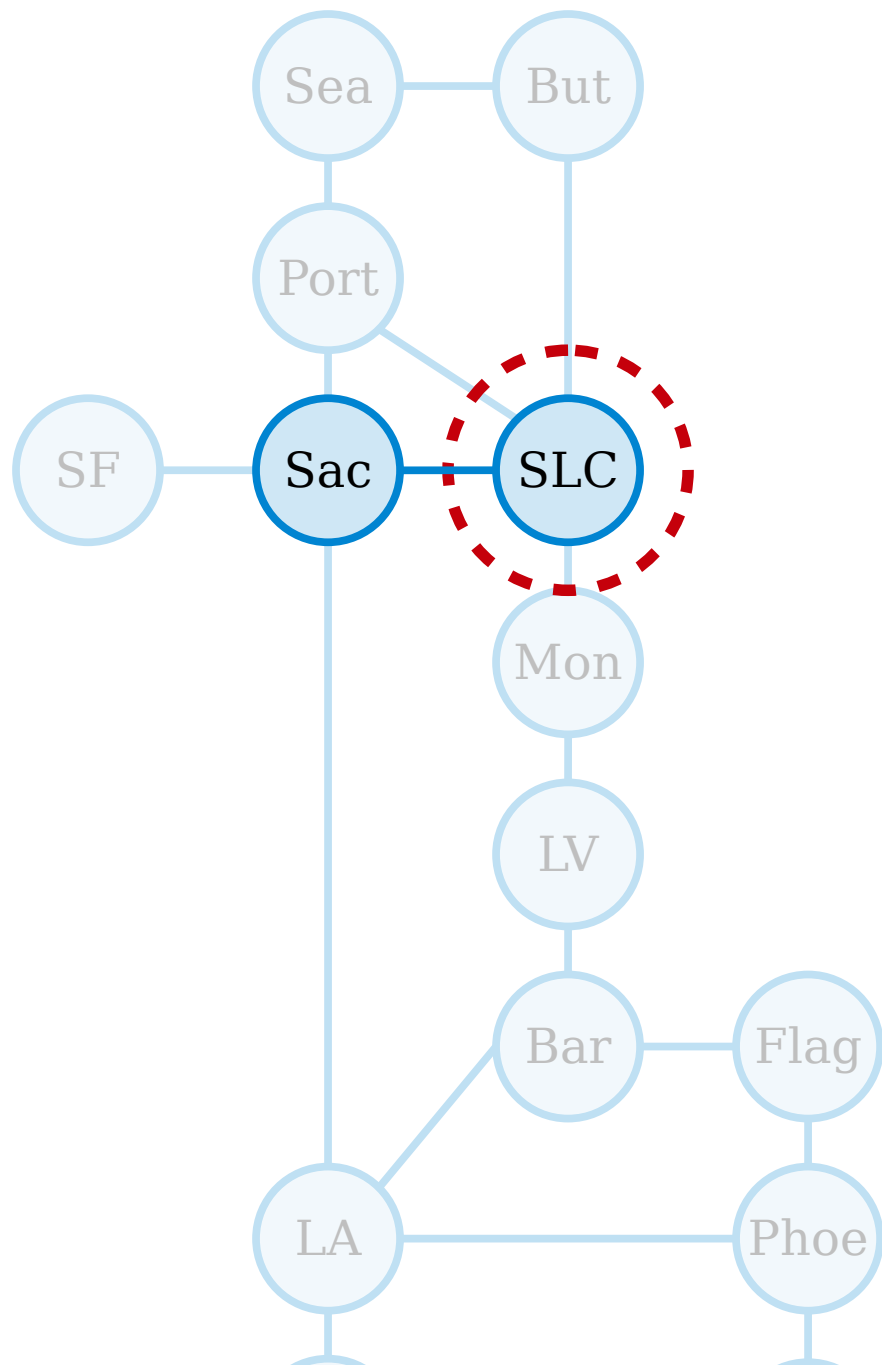


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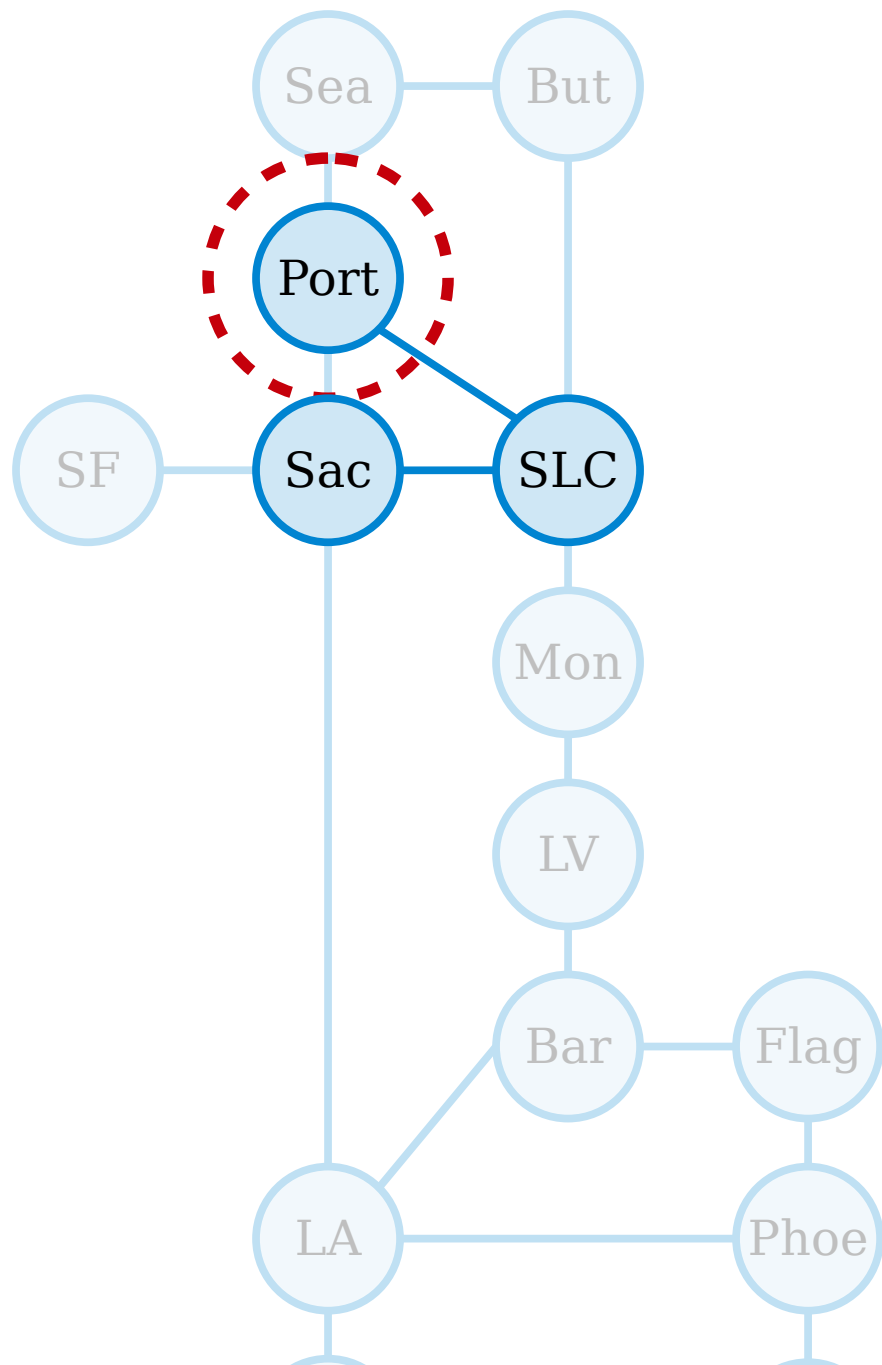
Sac, SLC

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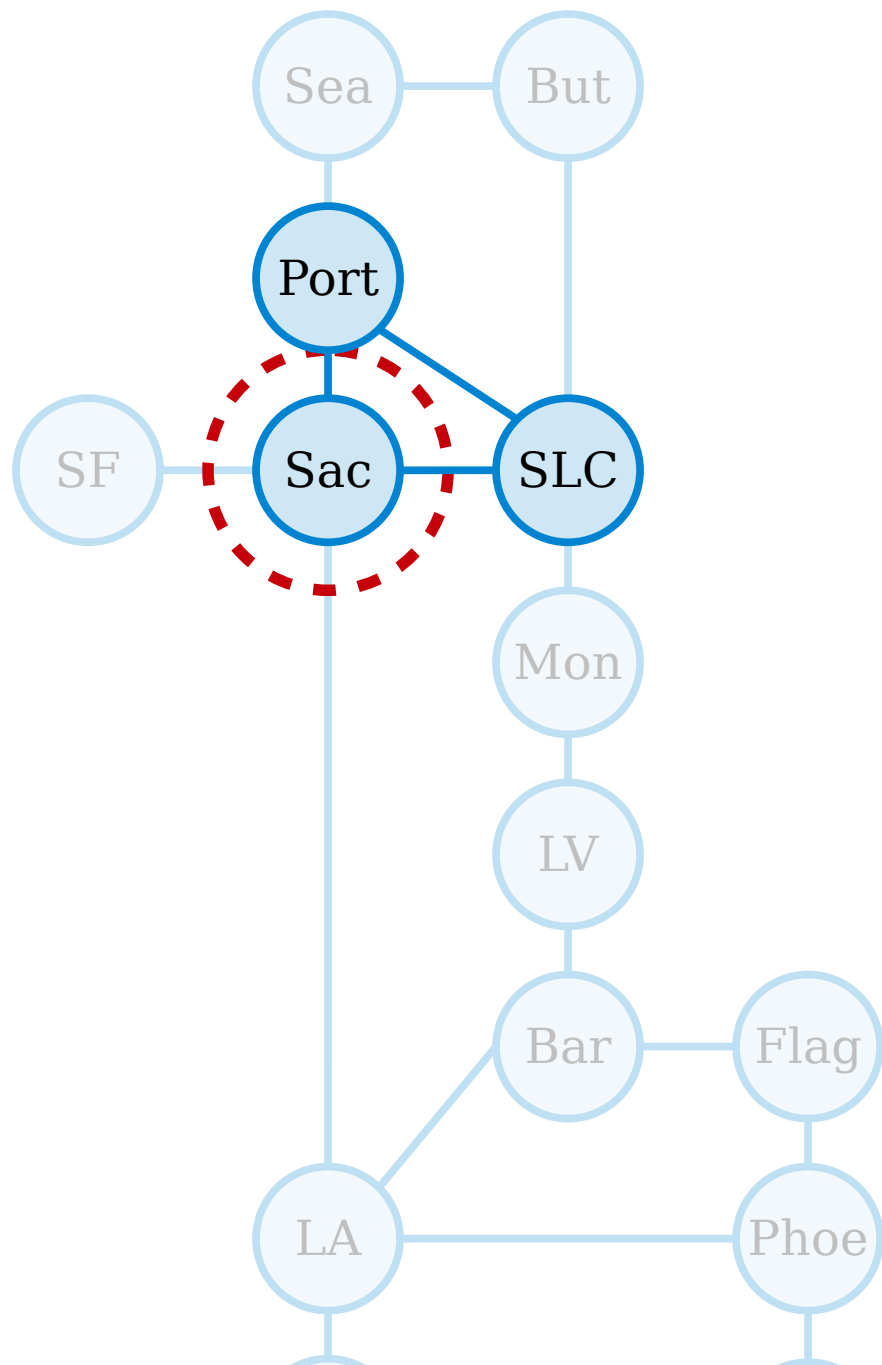
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Sac, SLC, Port

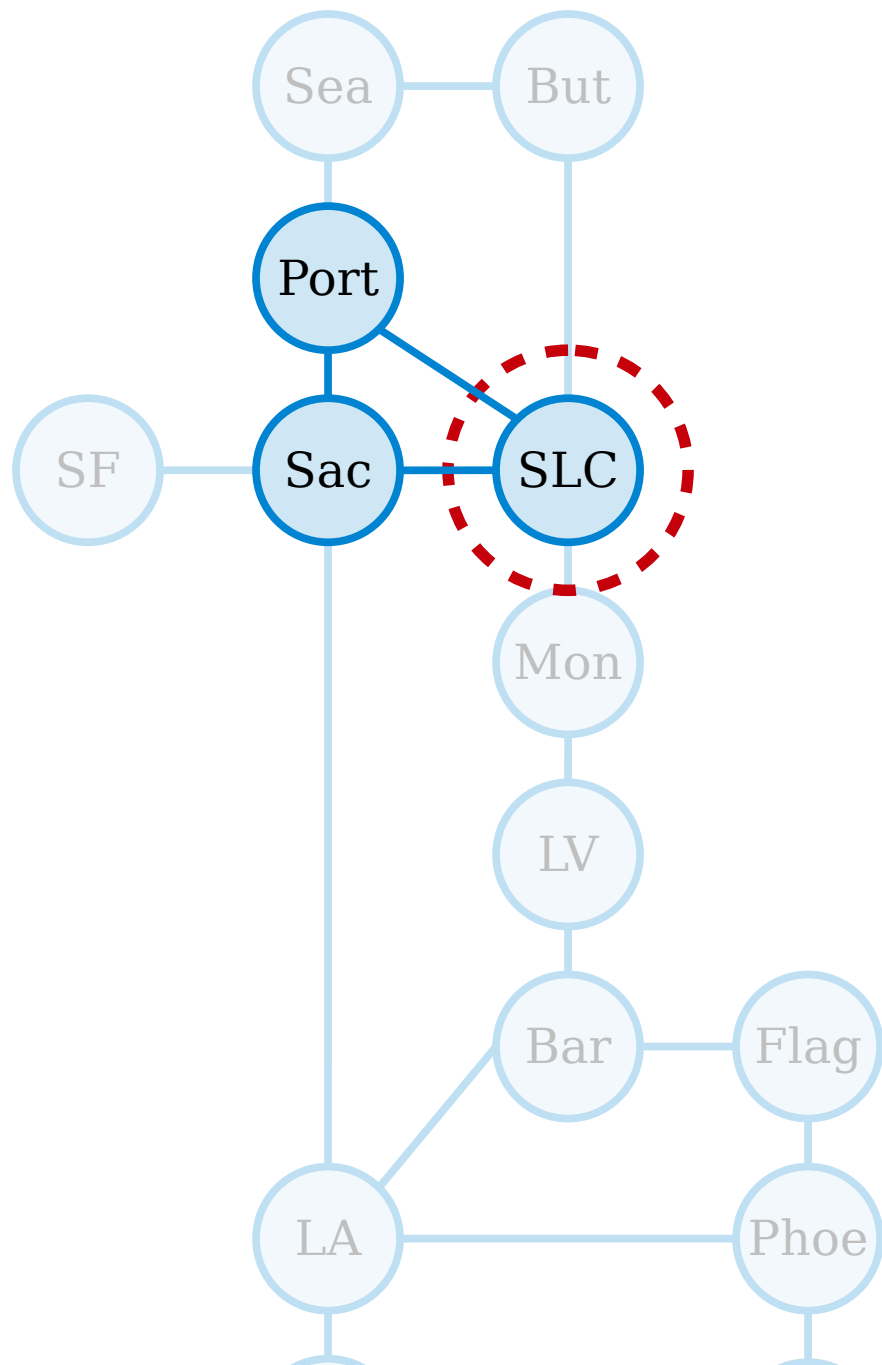


A **path** in a graph  $G = (V, E)$  is a sequence of one or more nodes  $v_1, v_2, v_3, \dots, v_n$  such that any two consecutive nodes in the sequence are adjacent.

The **length** of the path  $v_1, \dots, v_n$  is  $n - 1$ .

A **cycle** in a graph is a path from a node back to itself. (By convention, a cycle cannot have length zero.)

A **simple path** in a graph is path that does not repeat any nodes or edges.



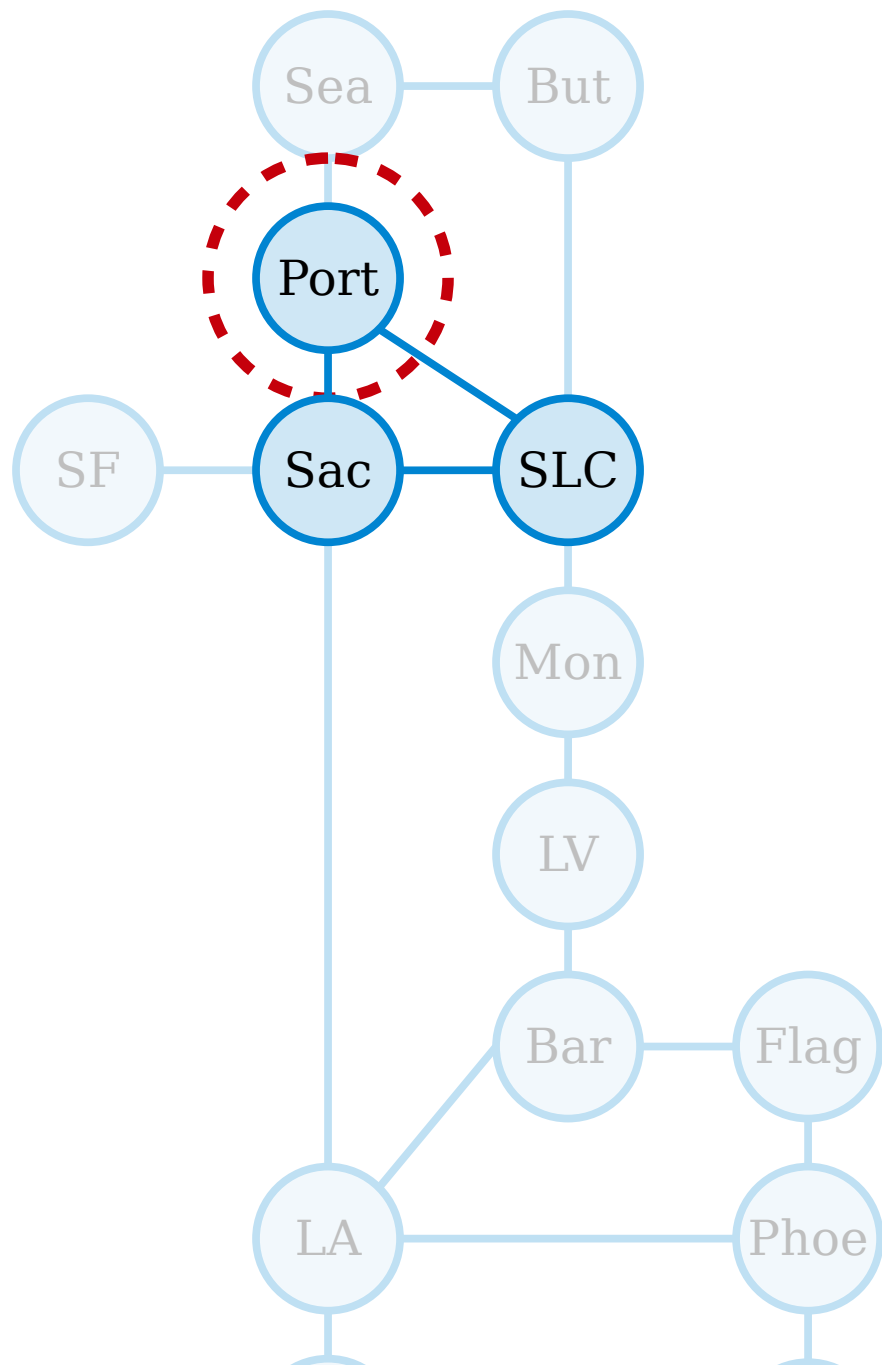
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Sac, SLC, Port, Sac, SLC



Sac, SLC, Port, Sac, SLC, Port

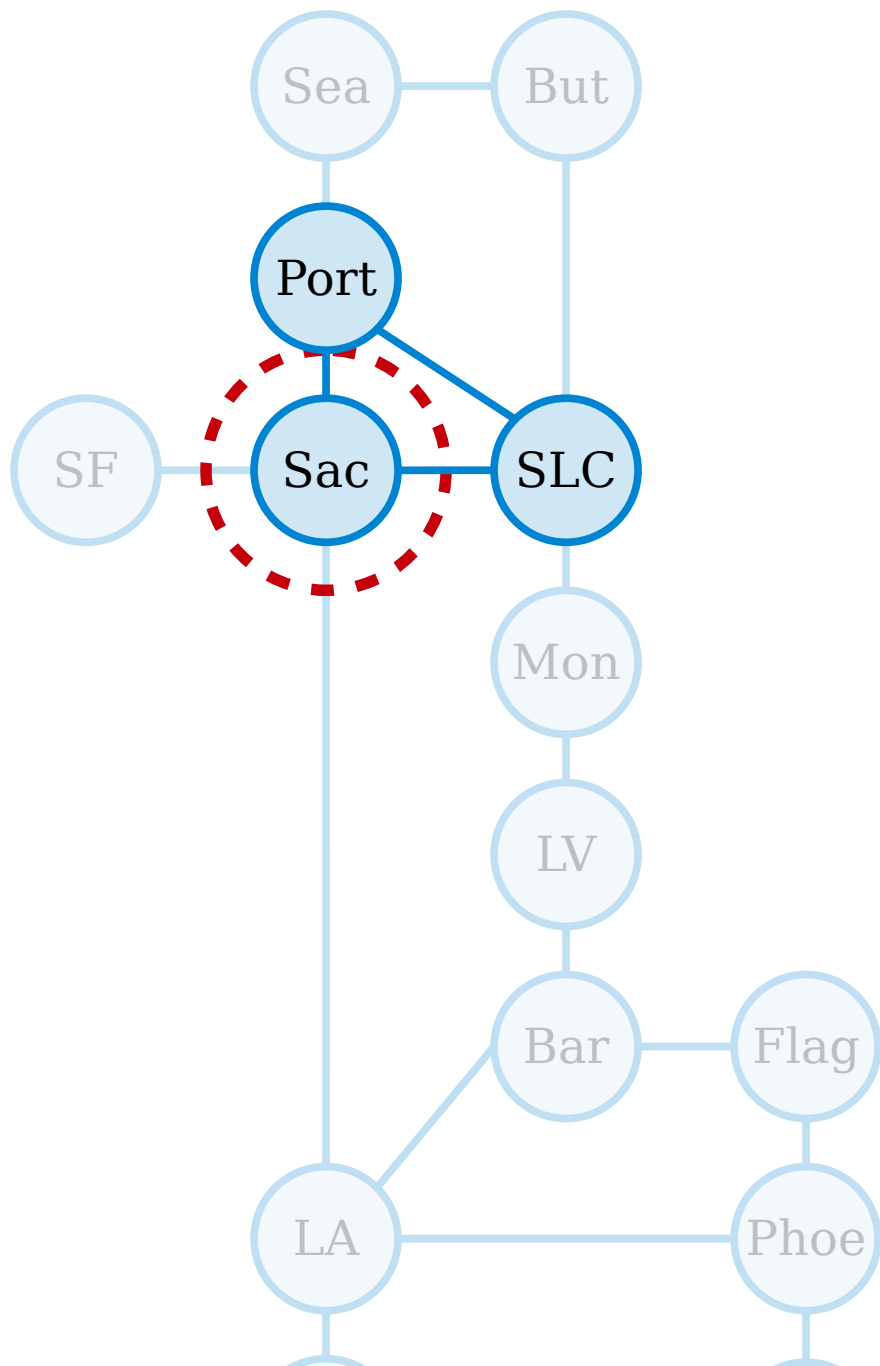
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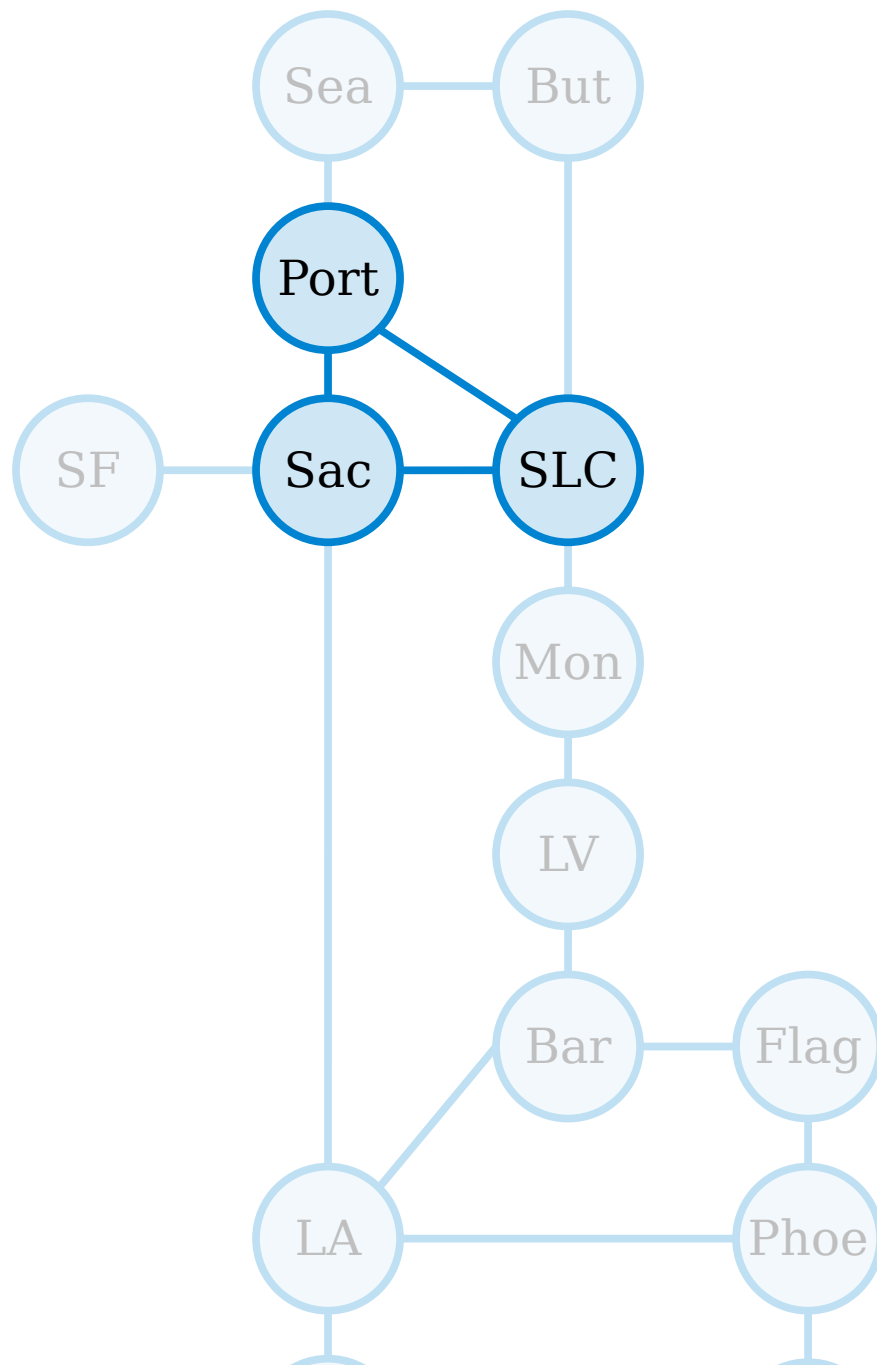
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Sac, SLC, Port, Sac, SLC, Port, Sac



Sac, SLC, Port, Sac, SLC, Port, Sac

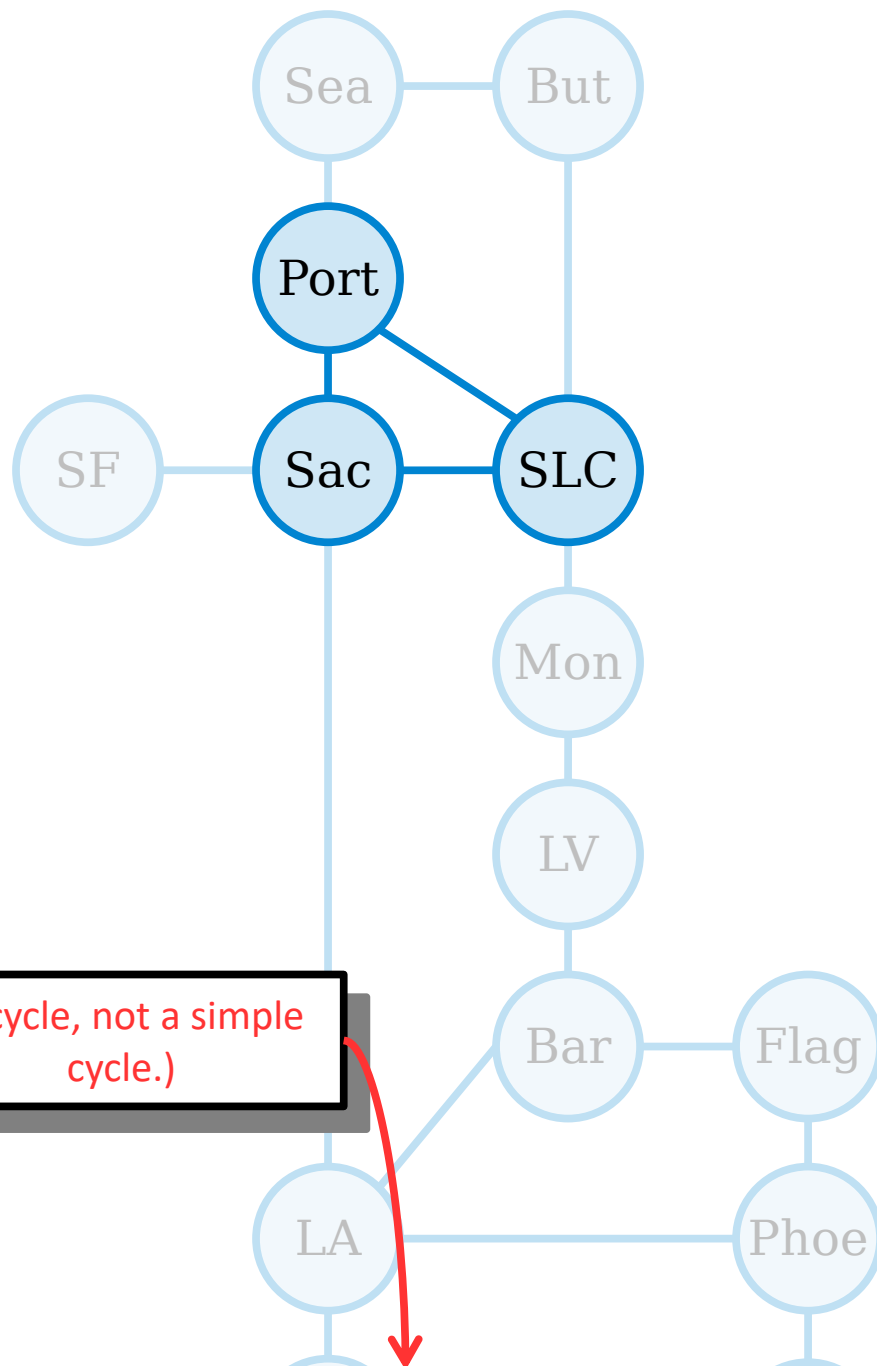
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A **simple cycle** in a graph is cycle that does not repeat any nodes or edges except the first/last node.



(A cycle, not a simple cycle.)

Sac, SLC, Port, Sac, SLC, Port, Sac

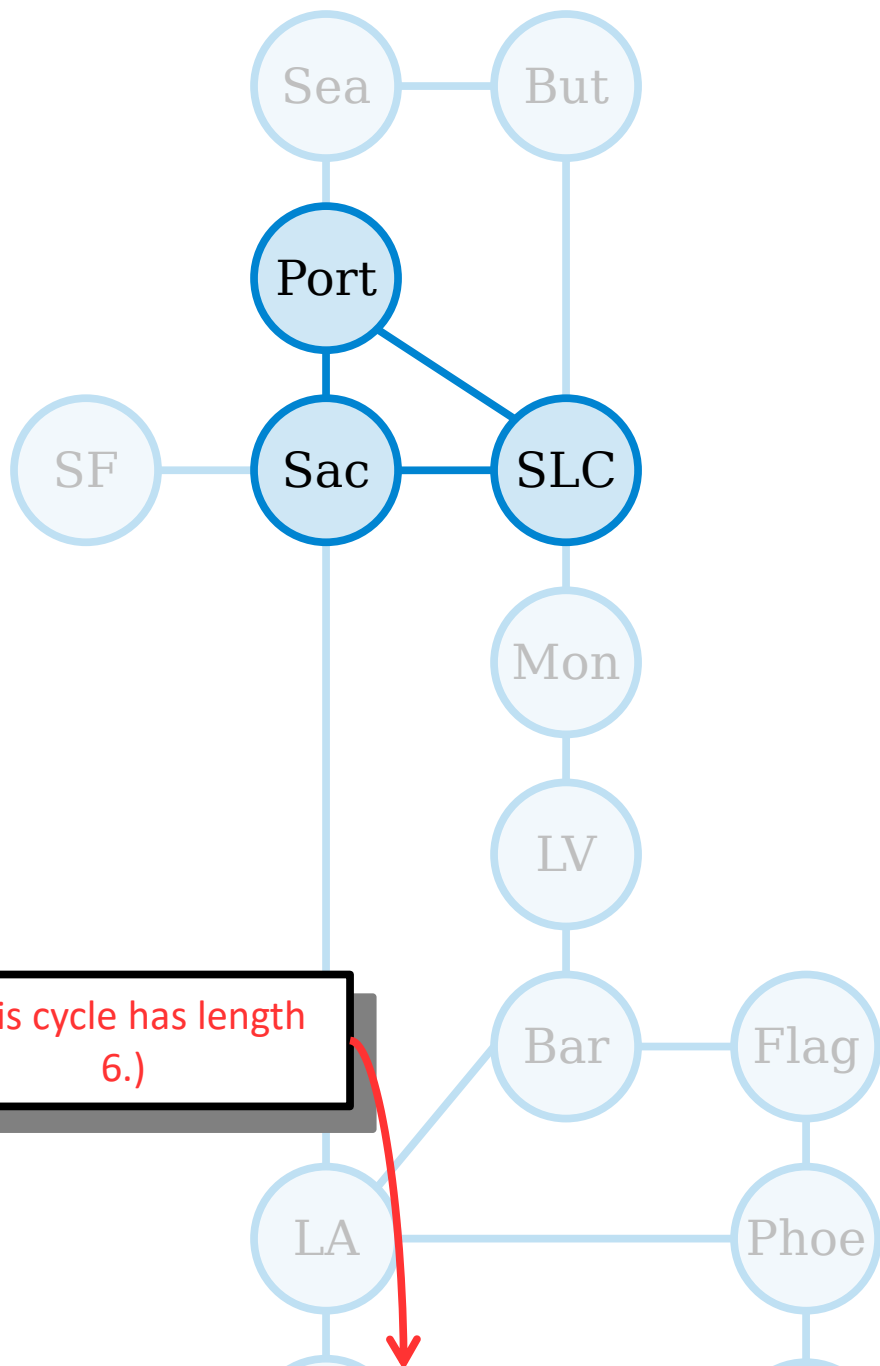
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(This cycle has length 6.)

Sac, SLC, Port, Sac, SLC, Port, Sac

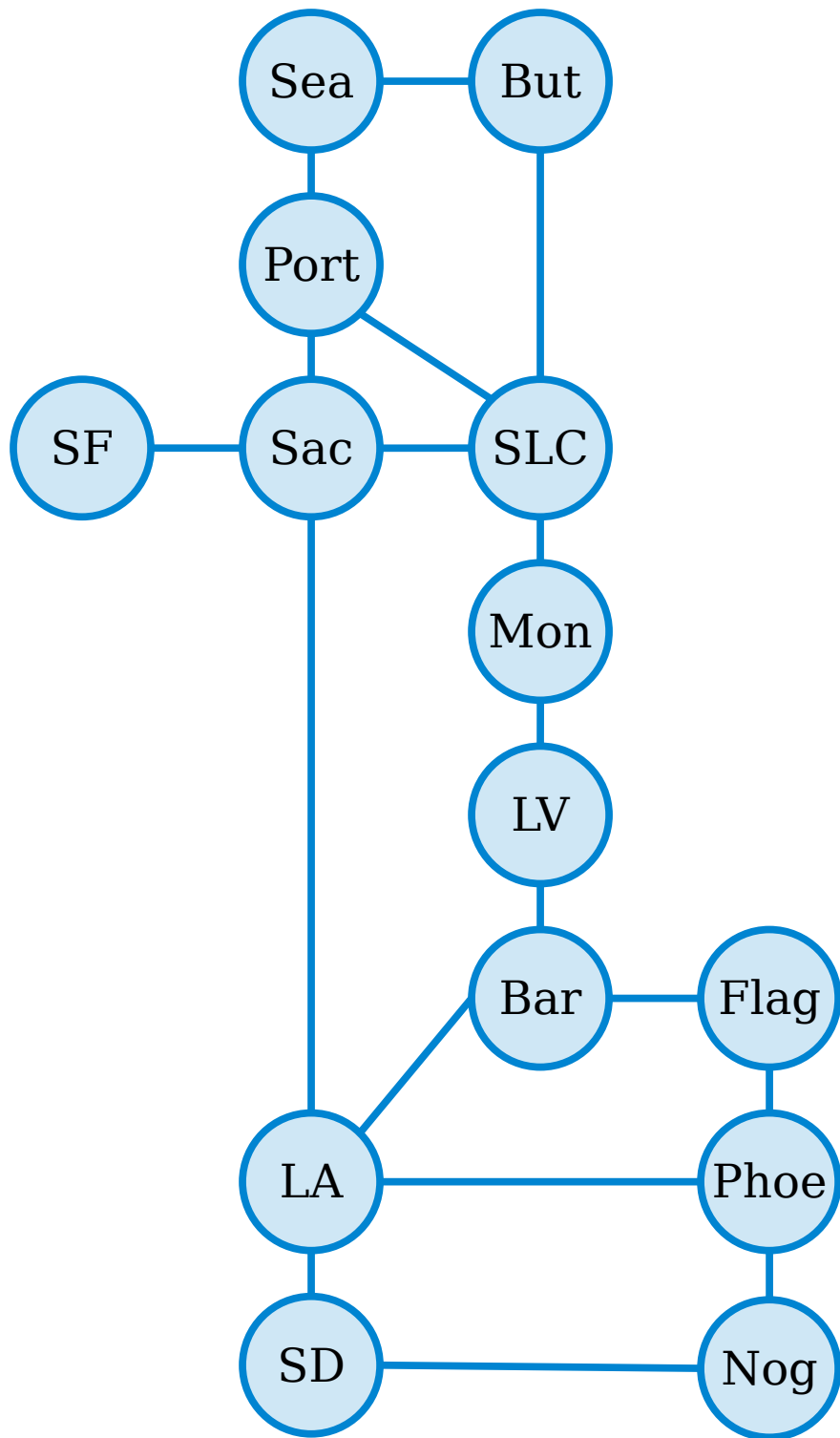
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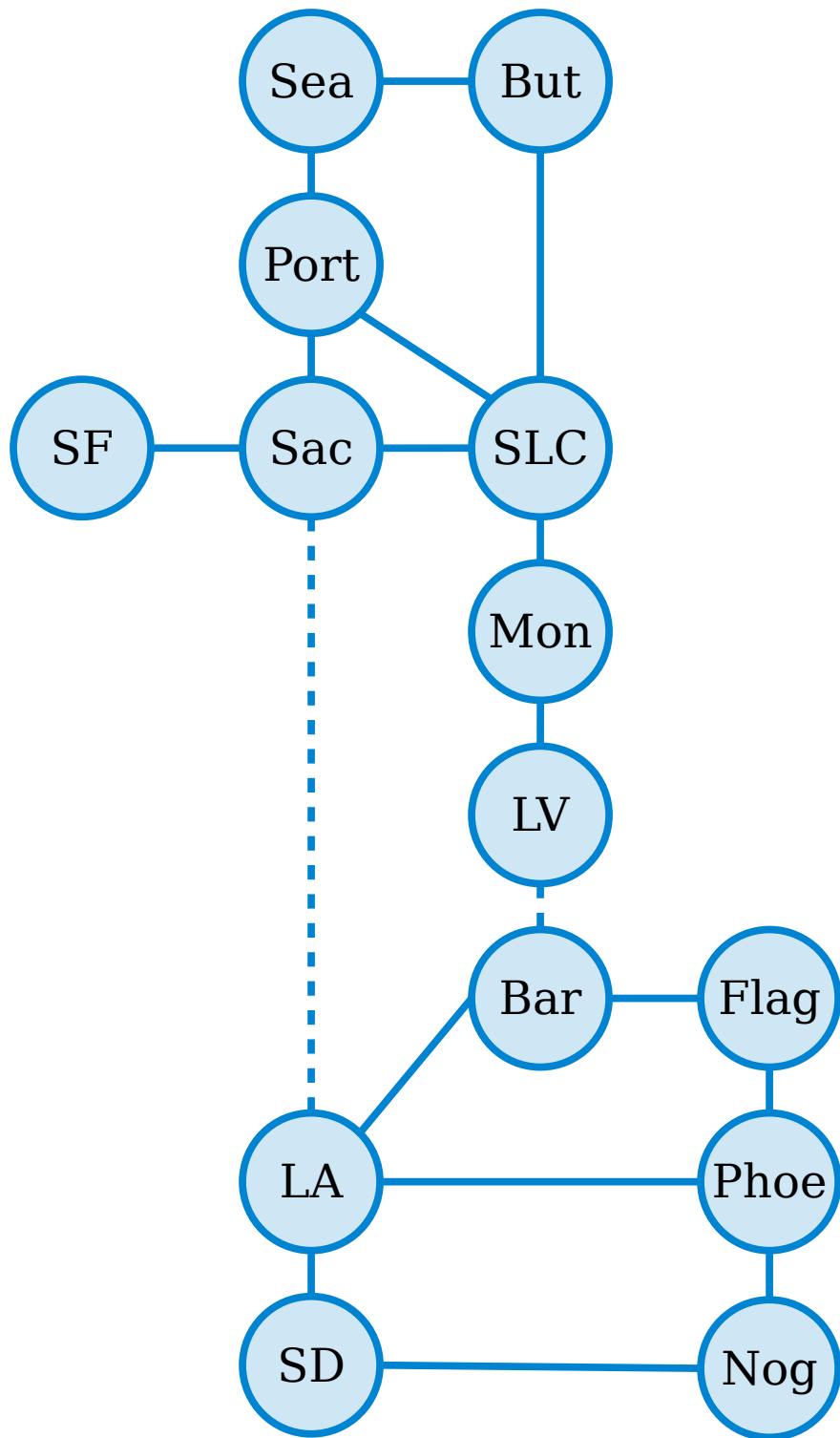
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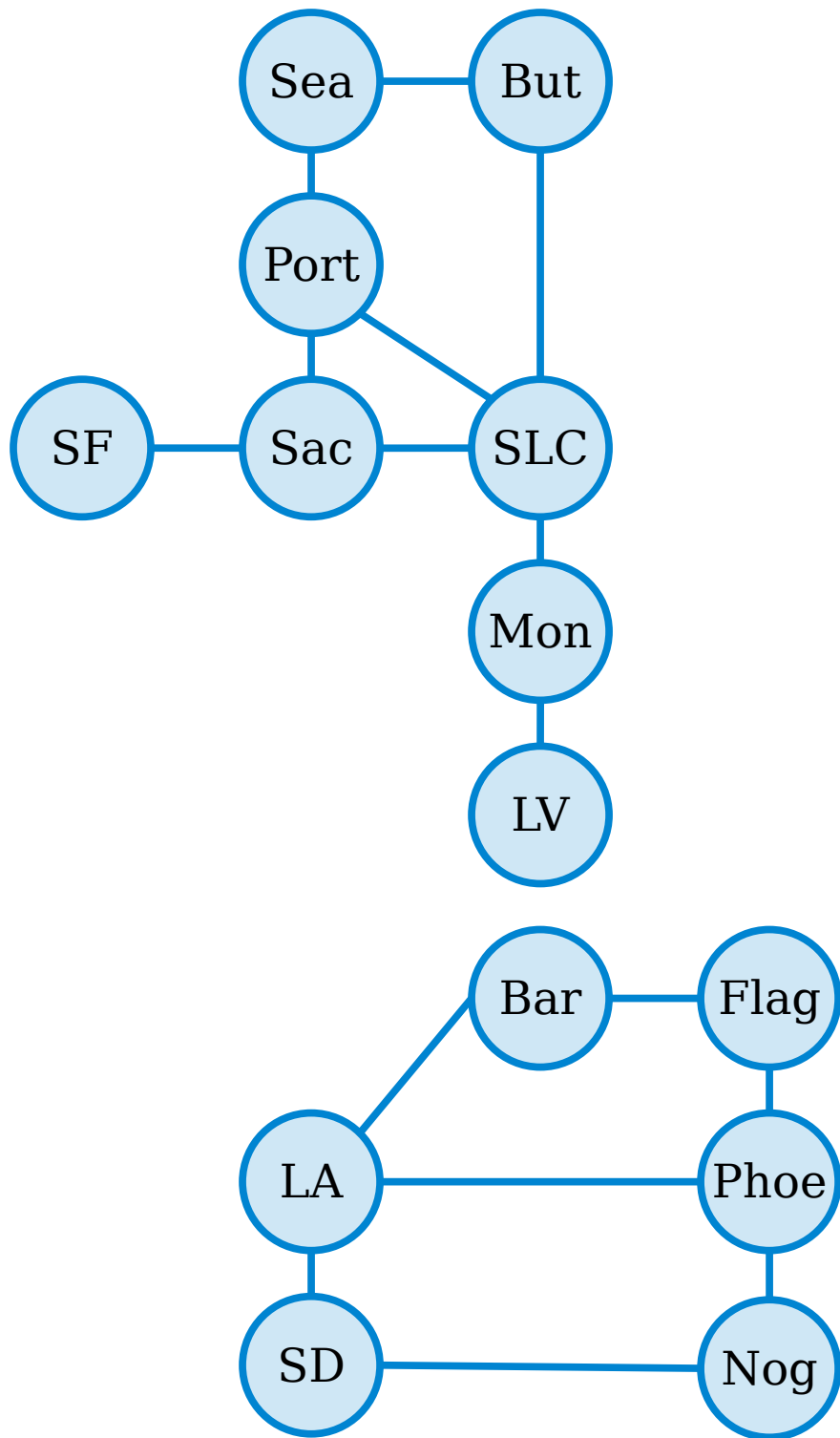
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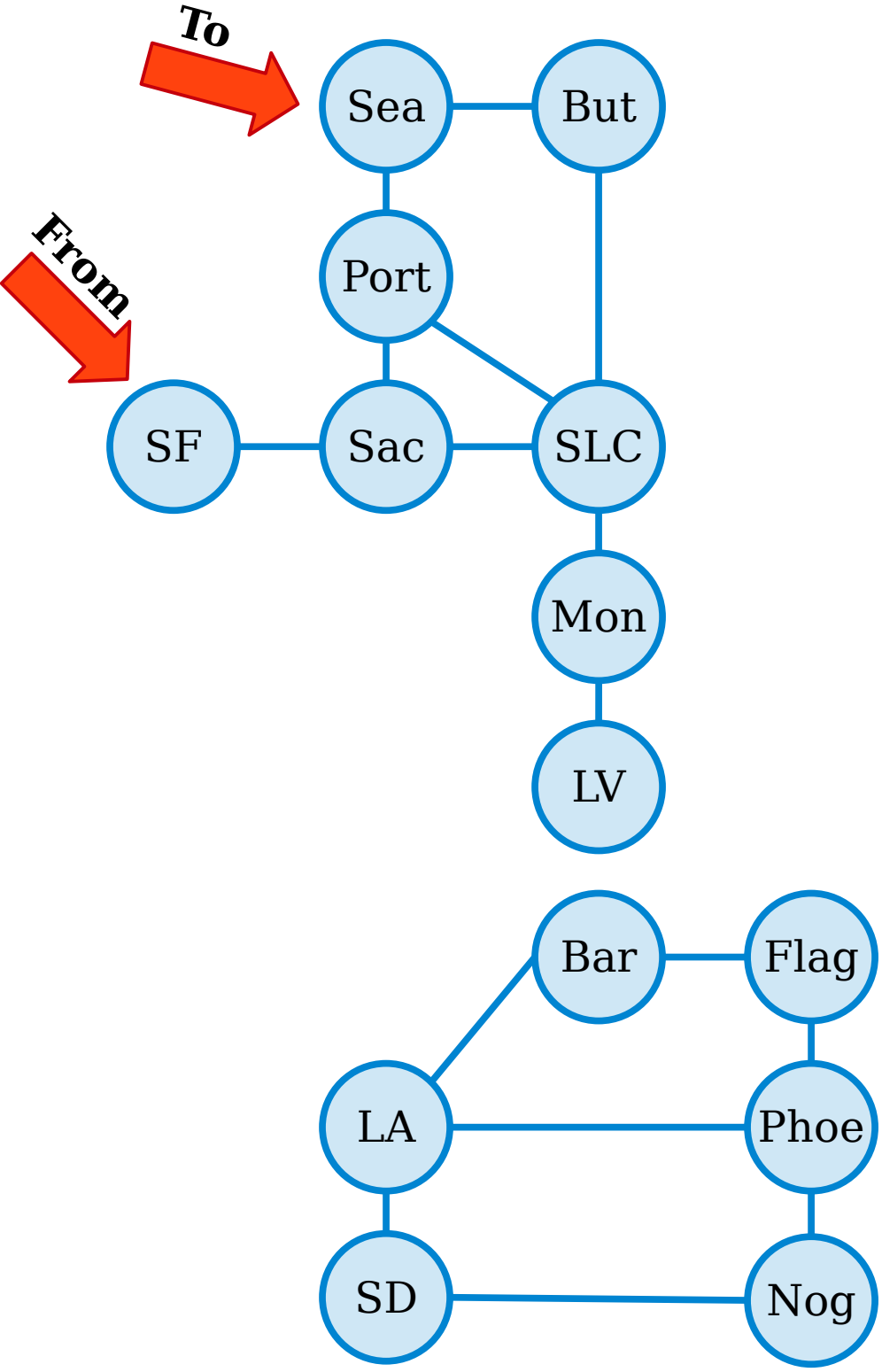


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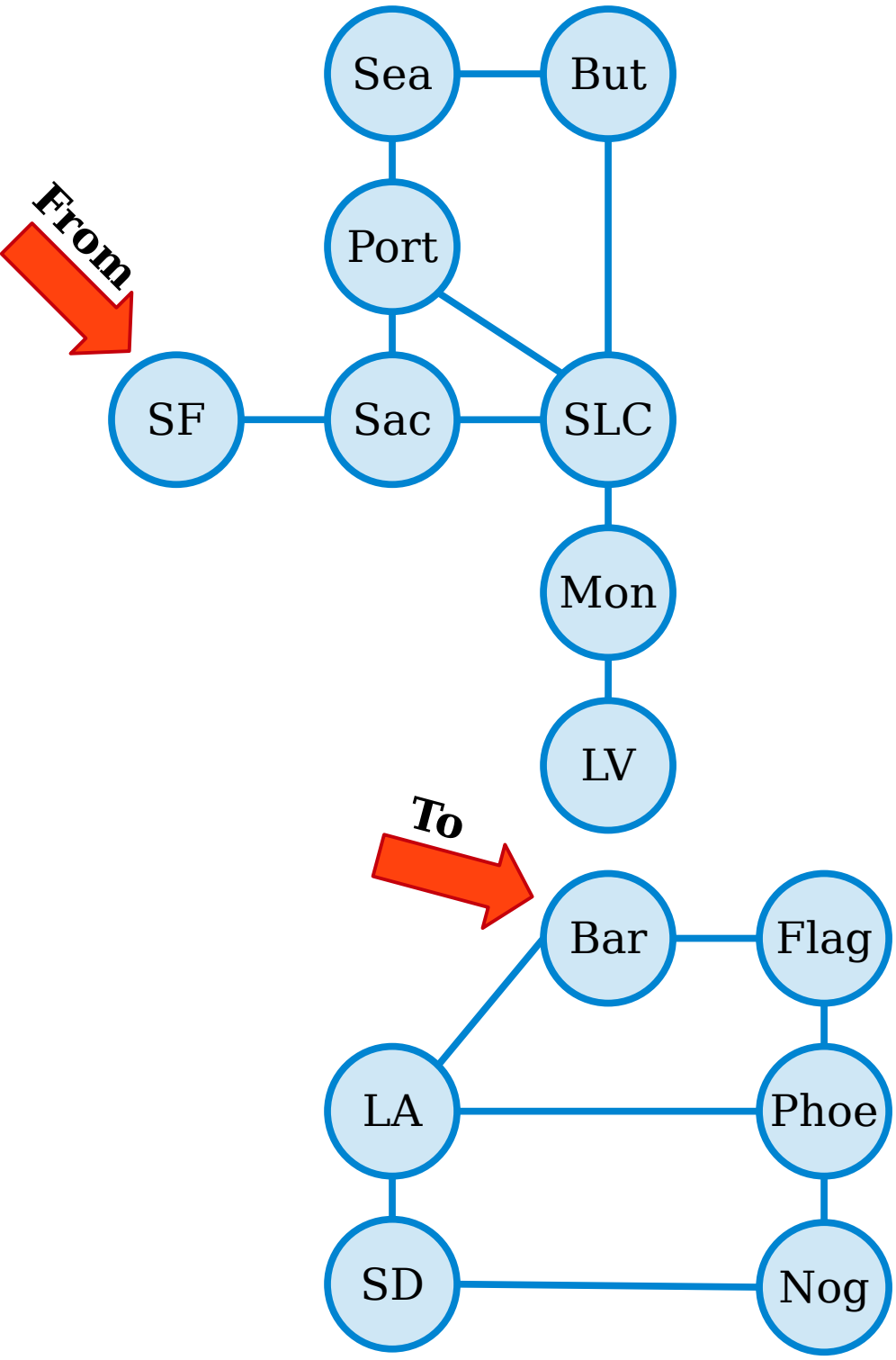
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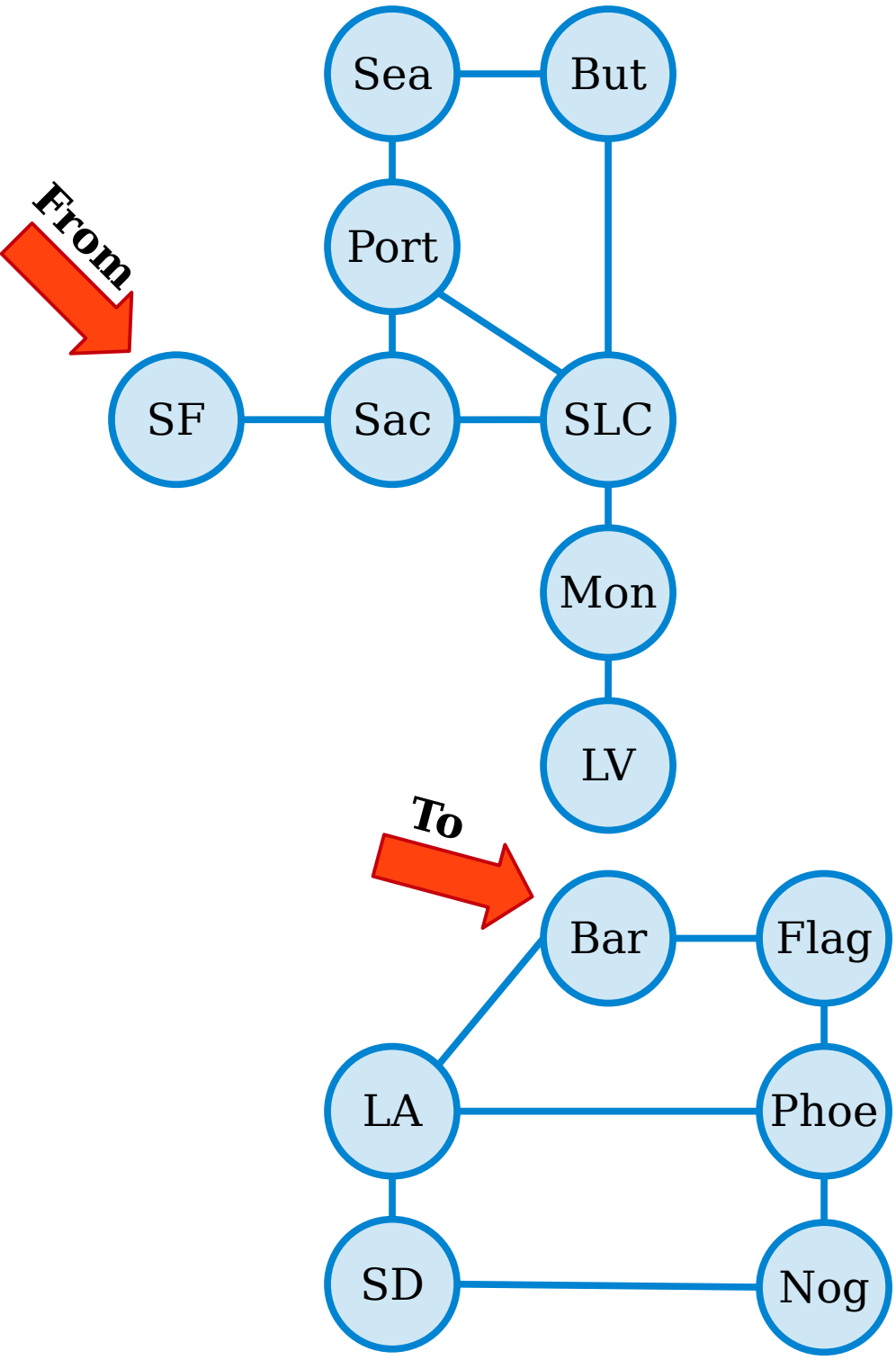
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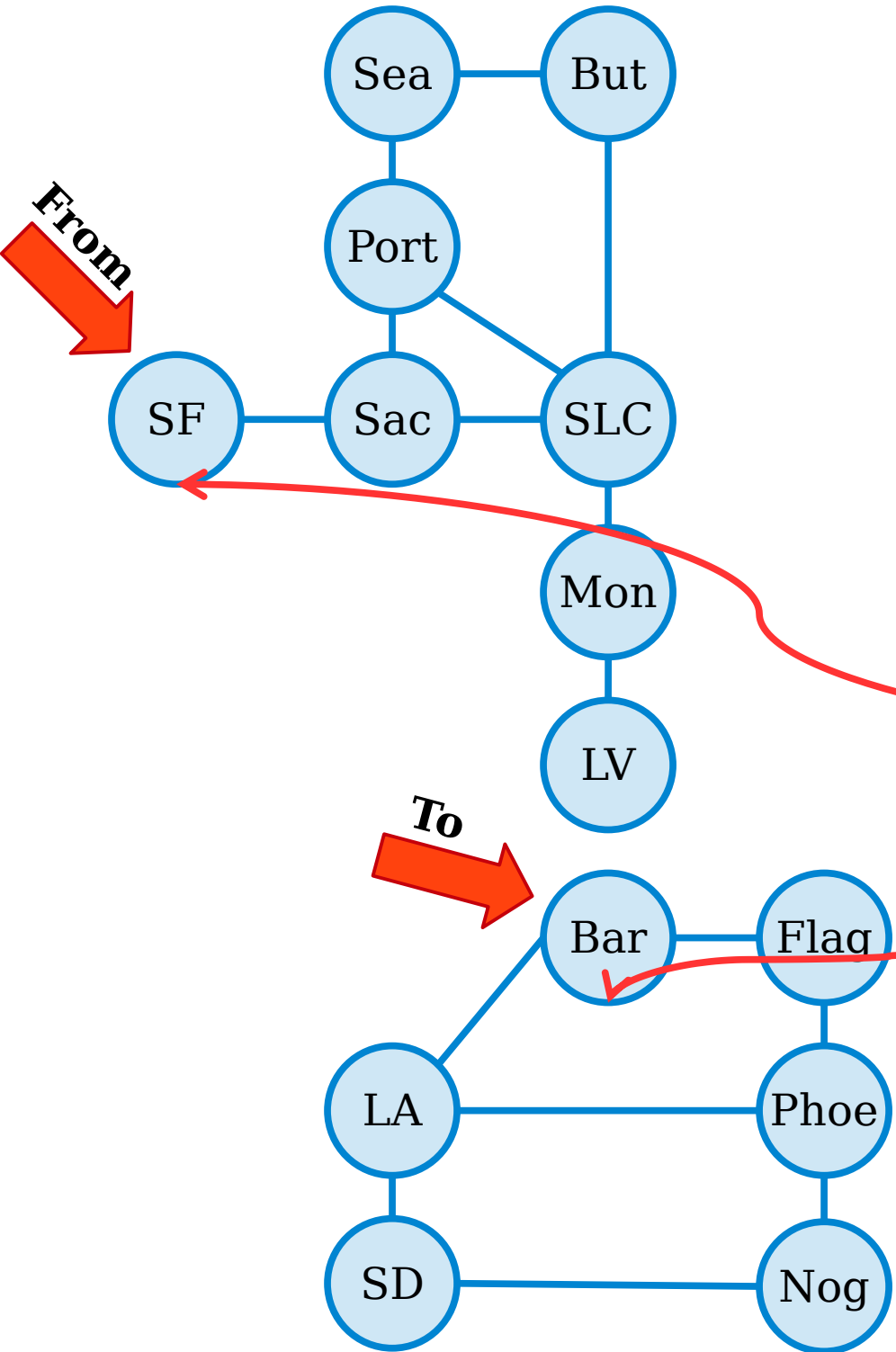
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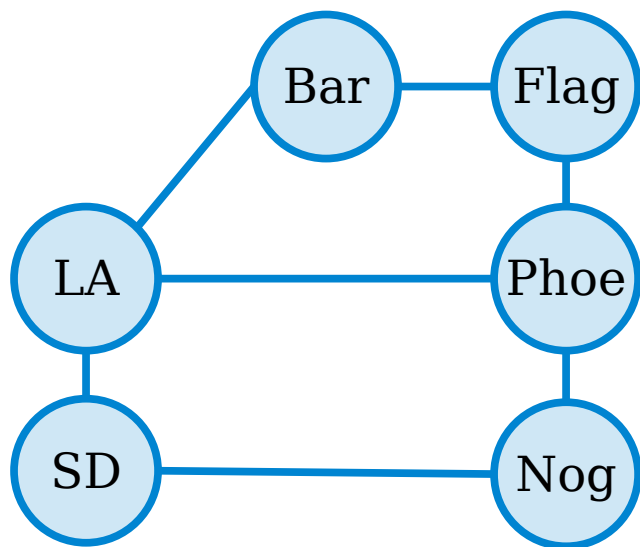
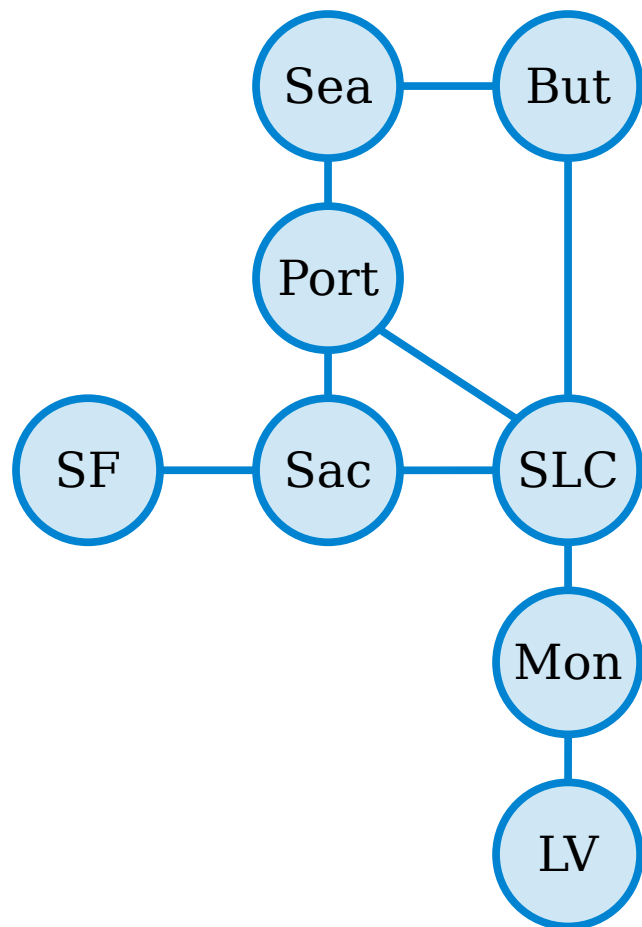
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(These nodes are not connected. No Grand Canyon for you.)

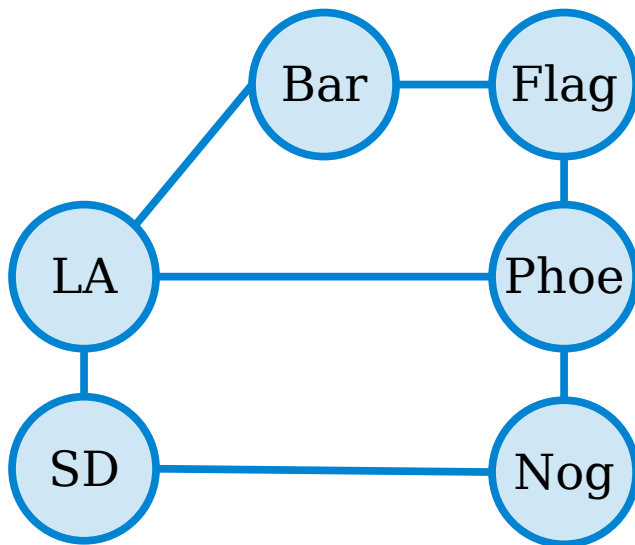
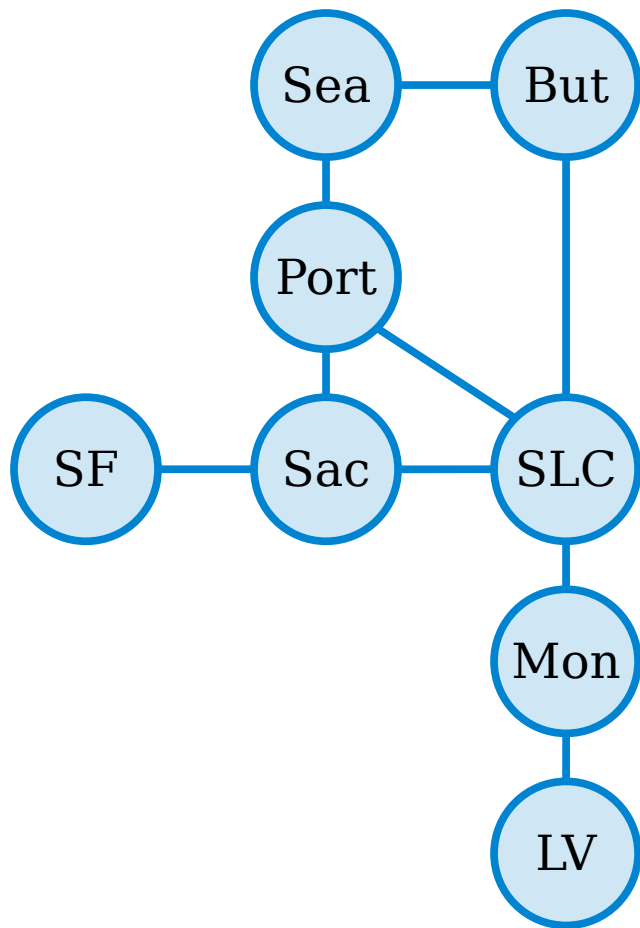




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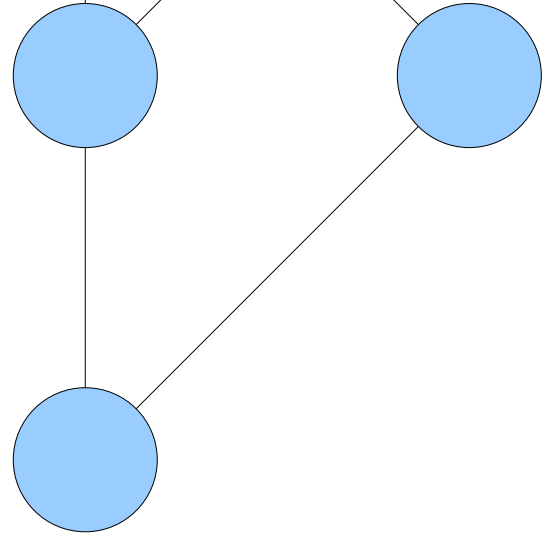
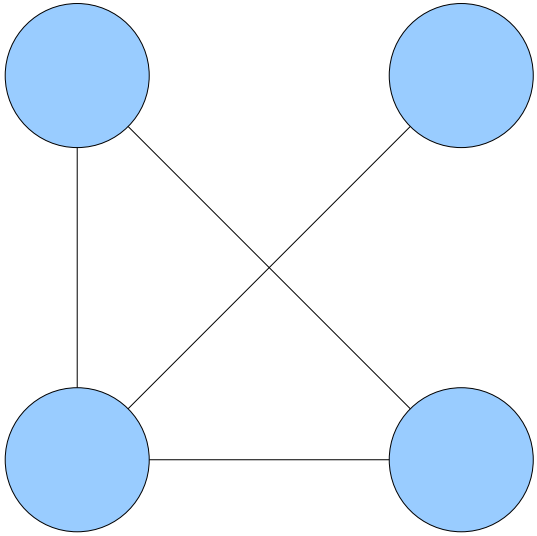
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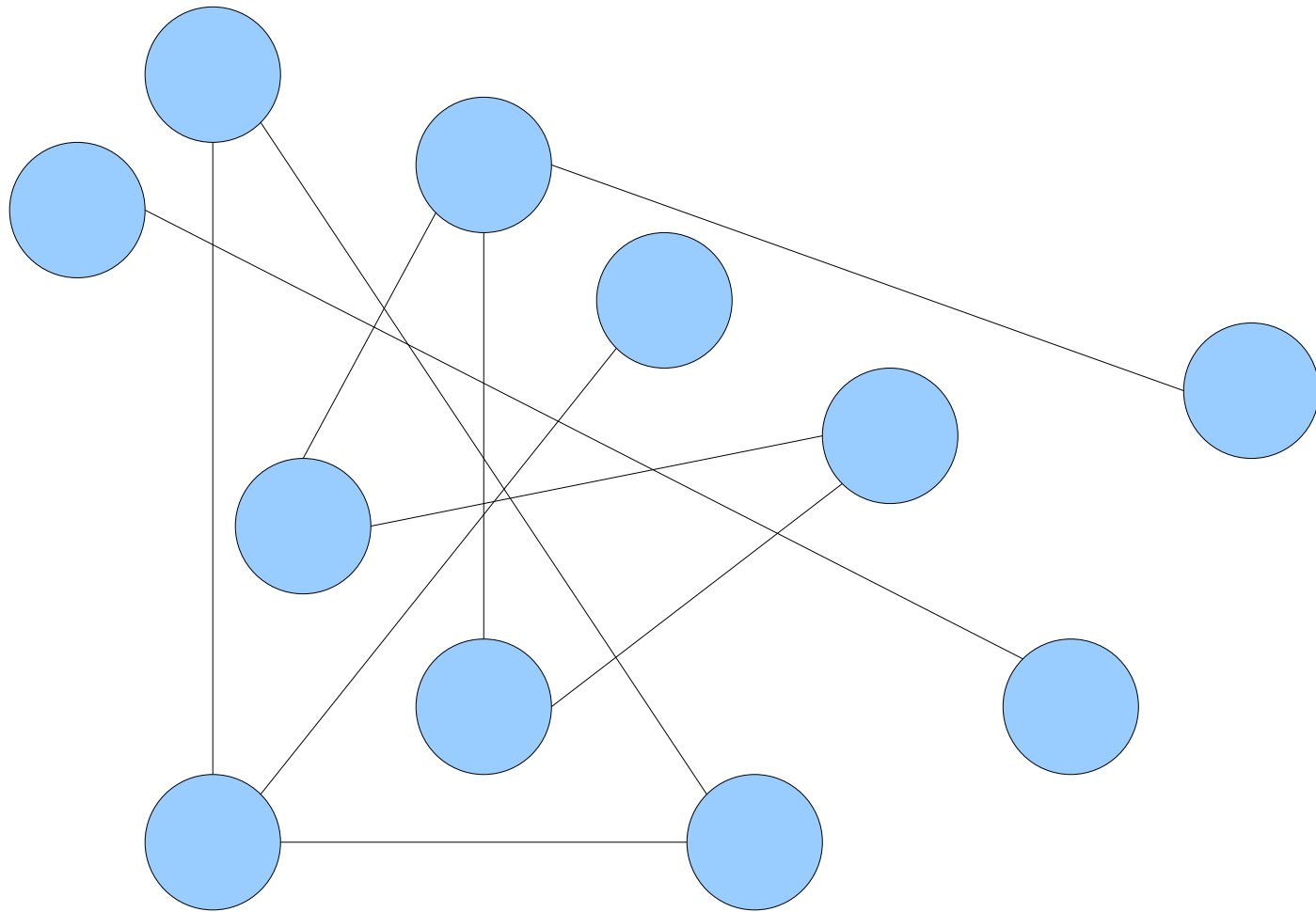
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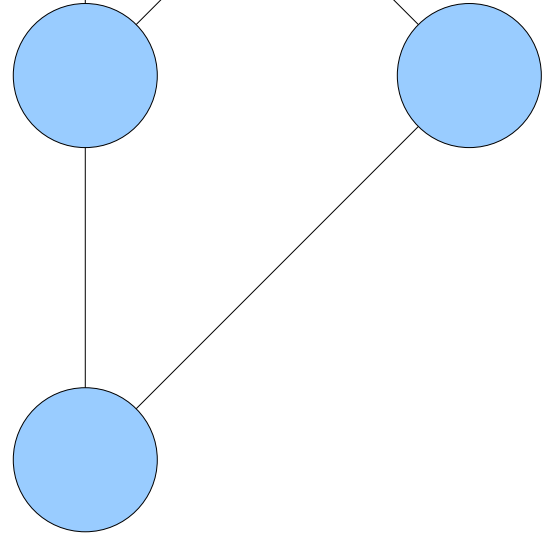
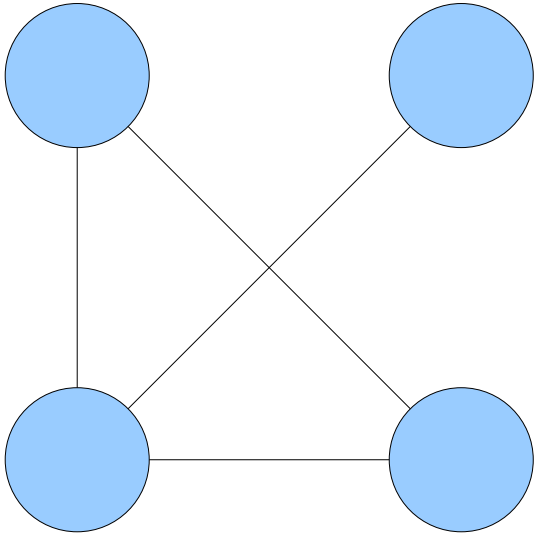
(This graph is not connected.)

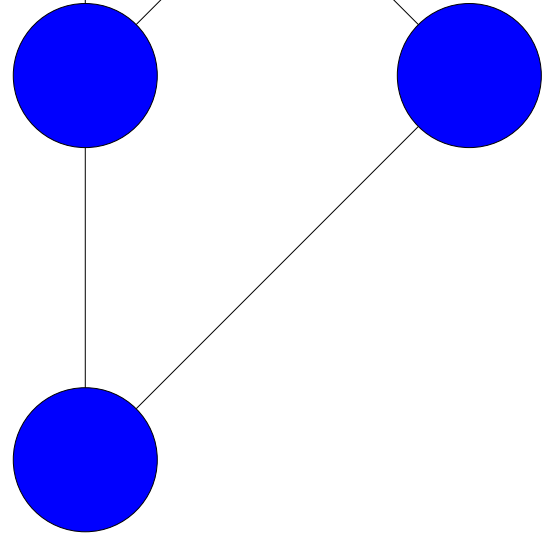
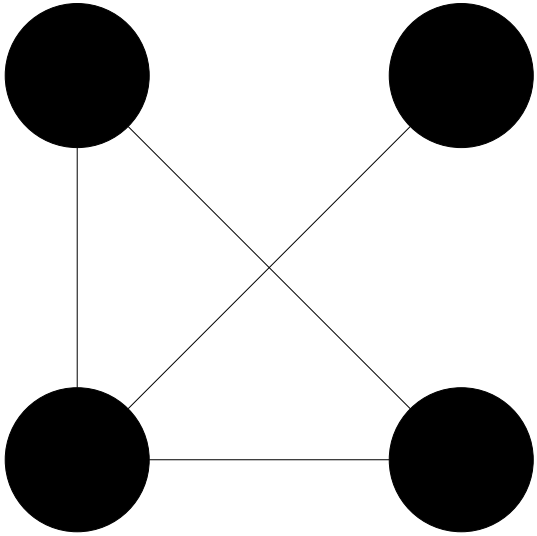
# Connected Components

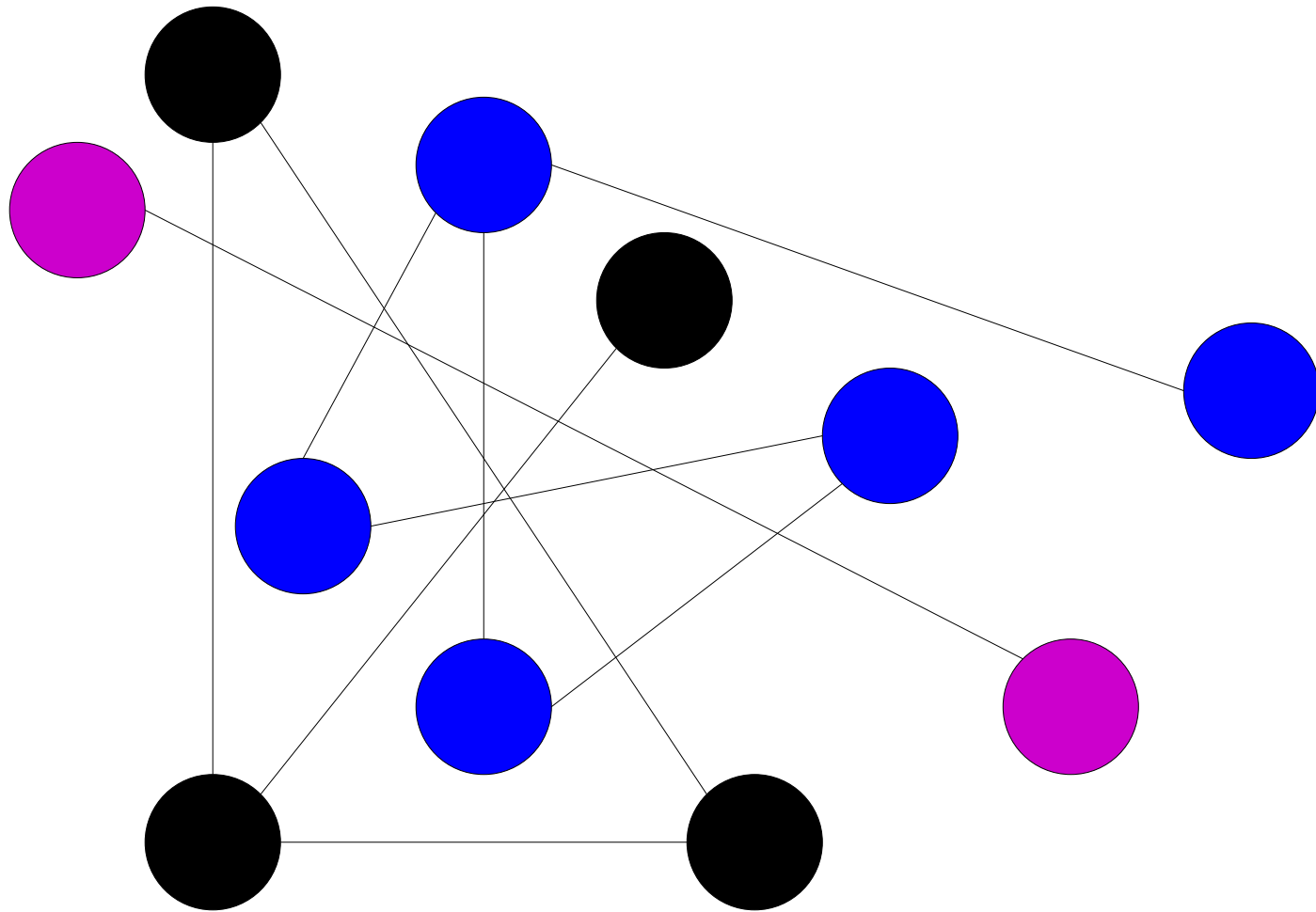


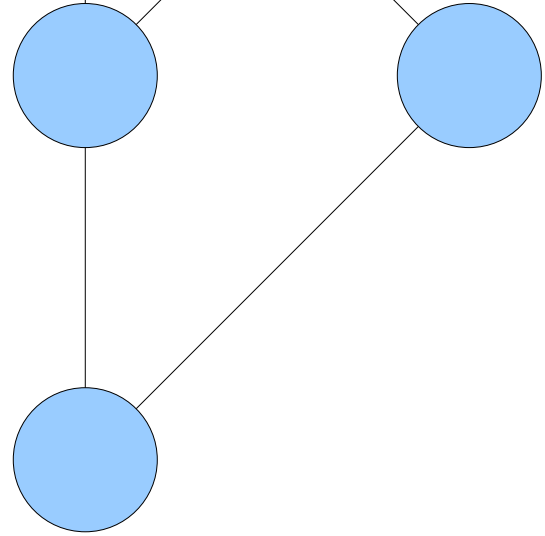
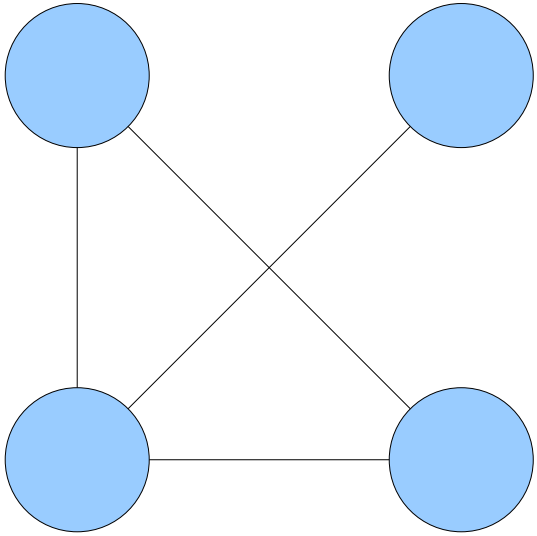


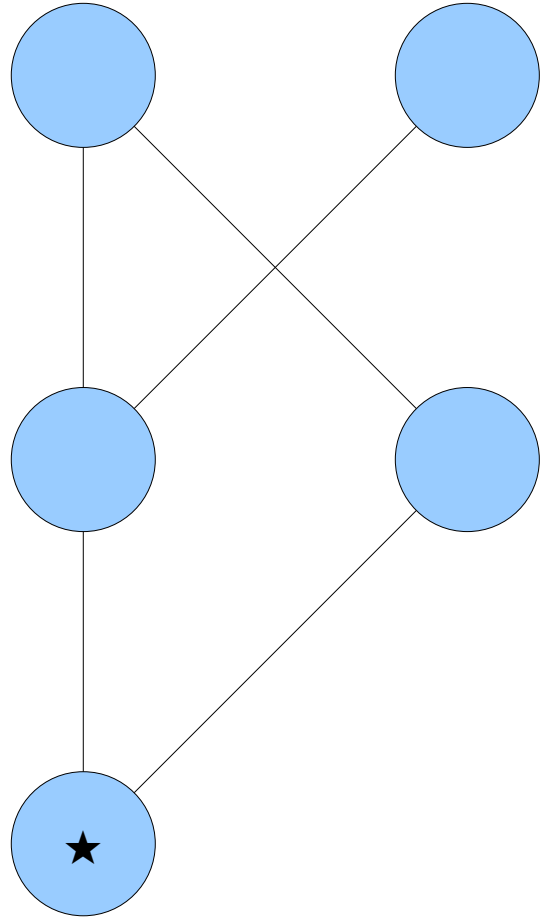
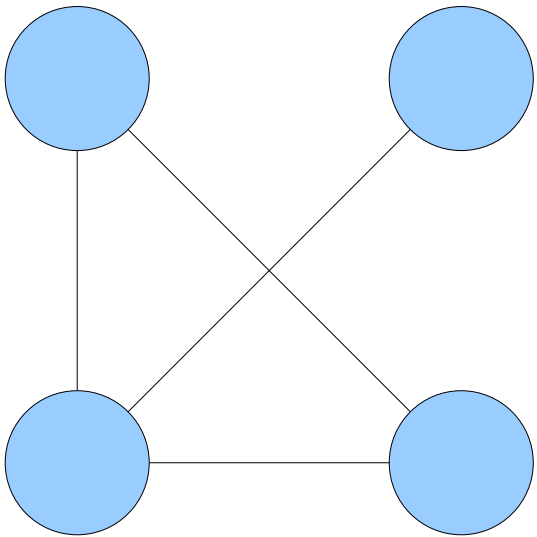


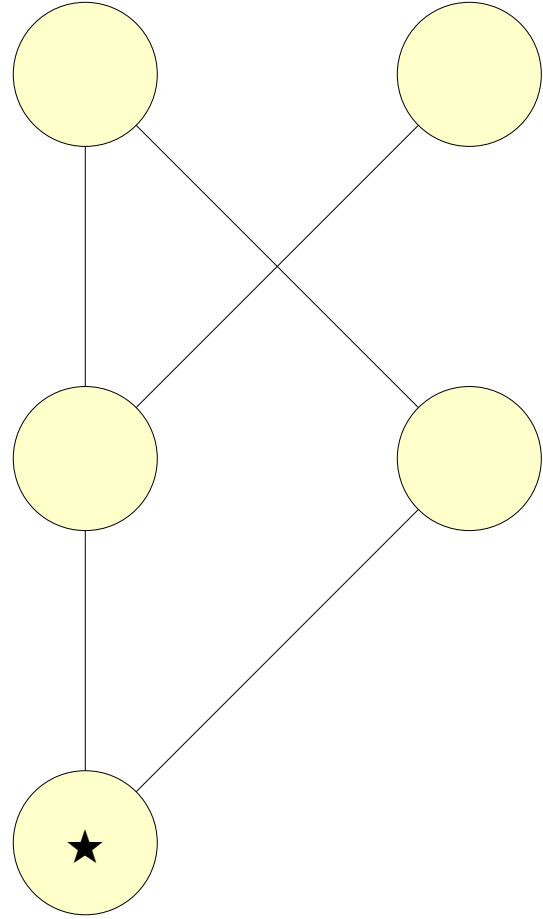
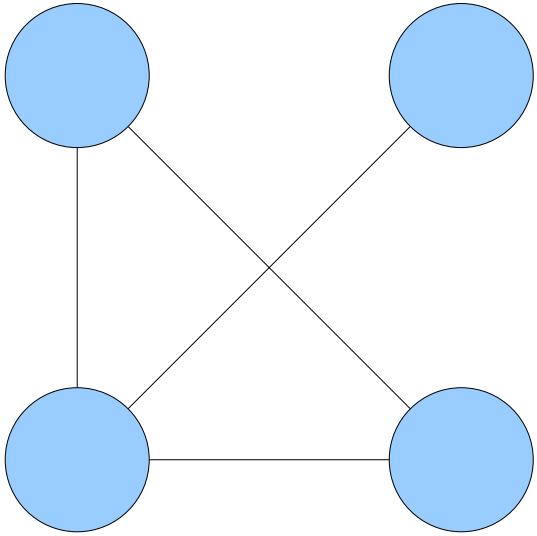












# Connected Components

Let  $G = (V, E)$  be a graph. For each  $v \in V$ , the **connected component** containing  $v$  is the set

$$[v] = \{ x \in V \mid v \text{ is connected to } x \}$$

Intuitively, a connected component is a “piece” of a graph in the sense we just talked about.

**Question:** How do we know that this particular definition of a “piece” of a graph is a good one?

**Goal:** Prove that any graph can be broken apart into different connected components.

We're trying to reason about some way of partitioning the nodes in a graph into different groups.

What structure have we studied that captures the idea of a partition?



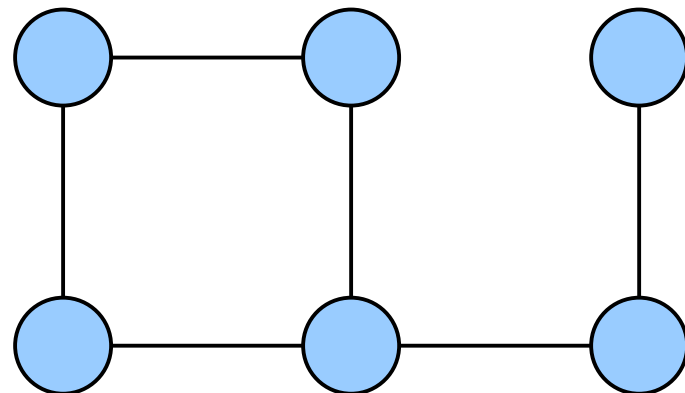
# Connectivity

**Claim:** For any graph  $G$ , the “is connected to” relation is an equivalence relation.

Is it reflexive?

Is it symmetric?

Is it transitive?



# Connectivity

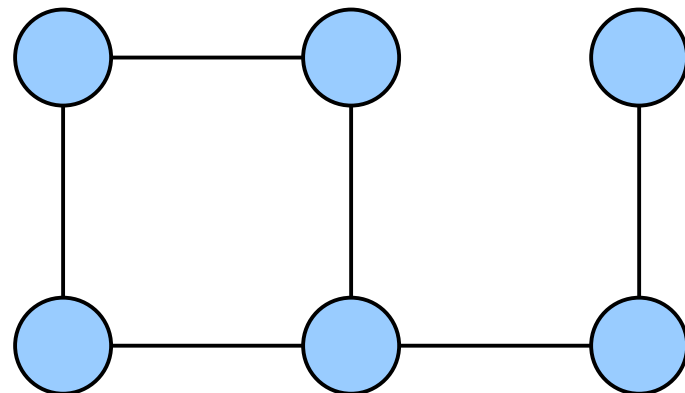
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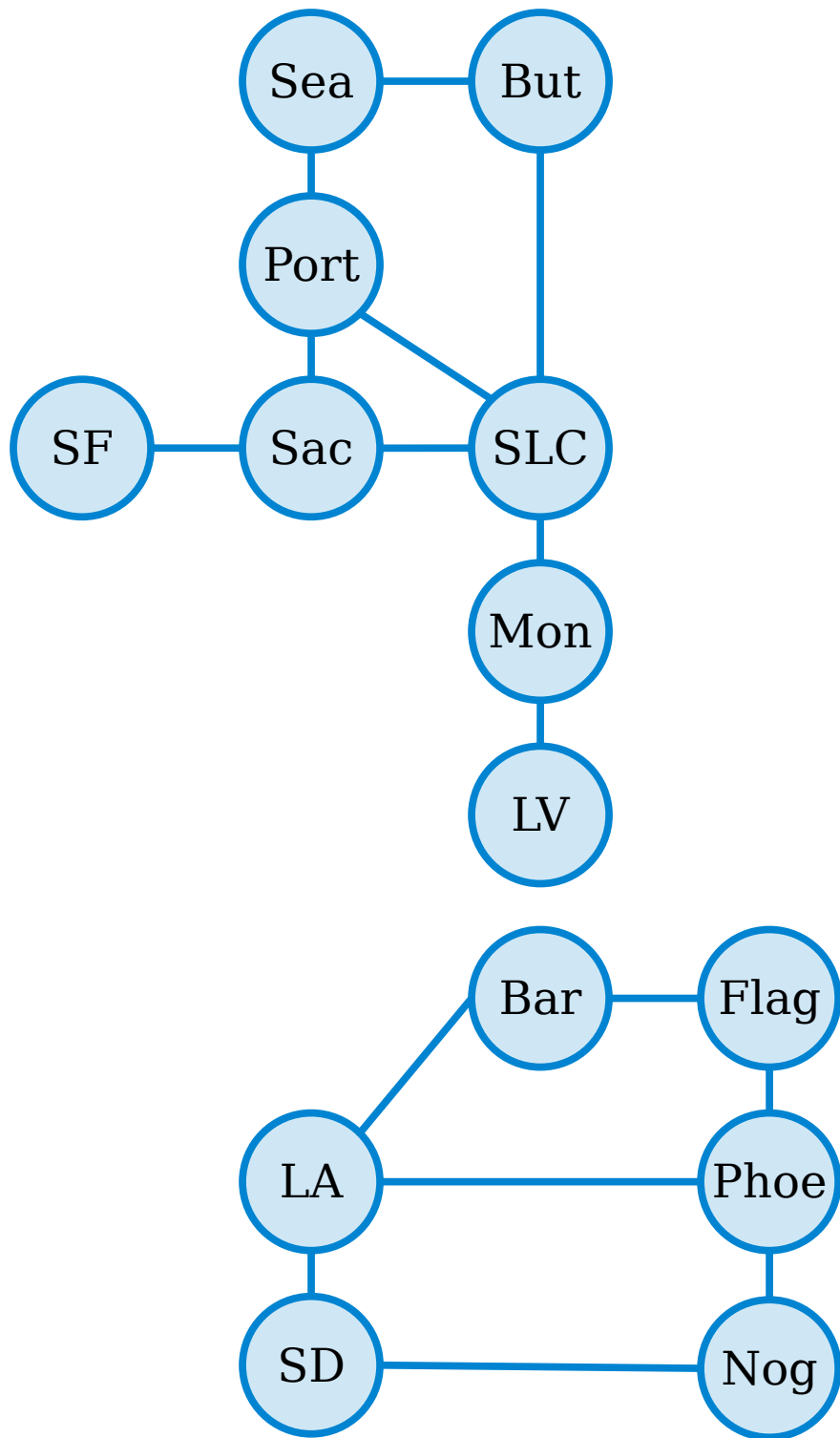
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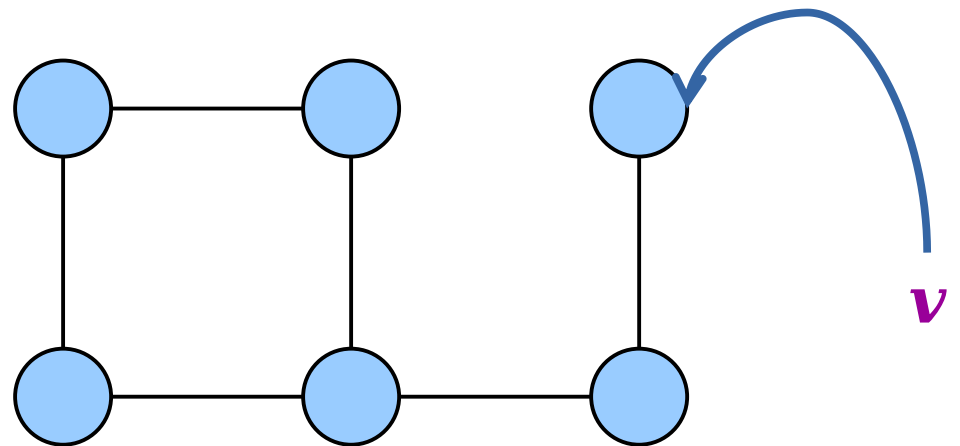
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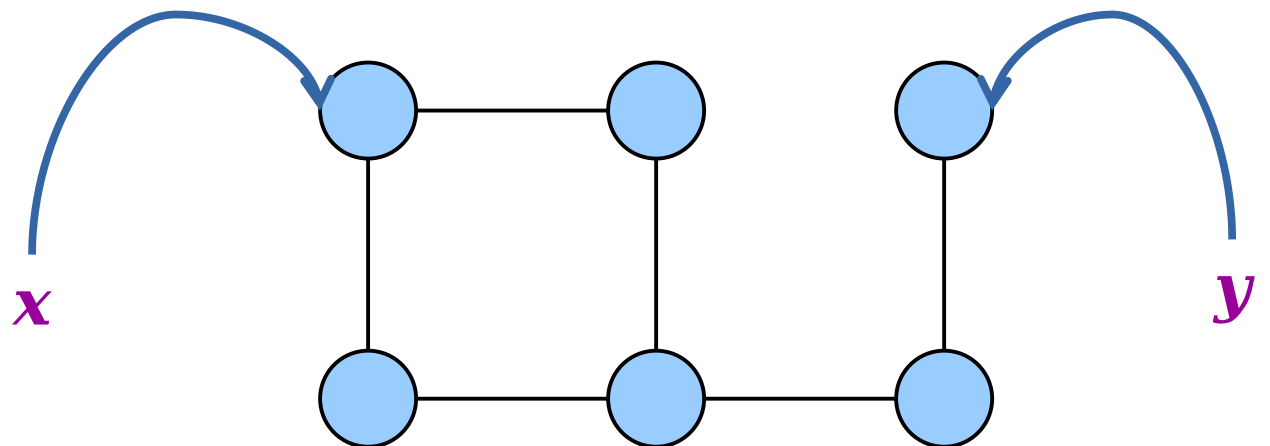
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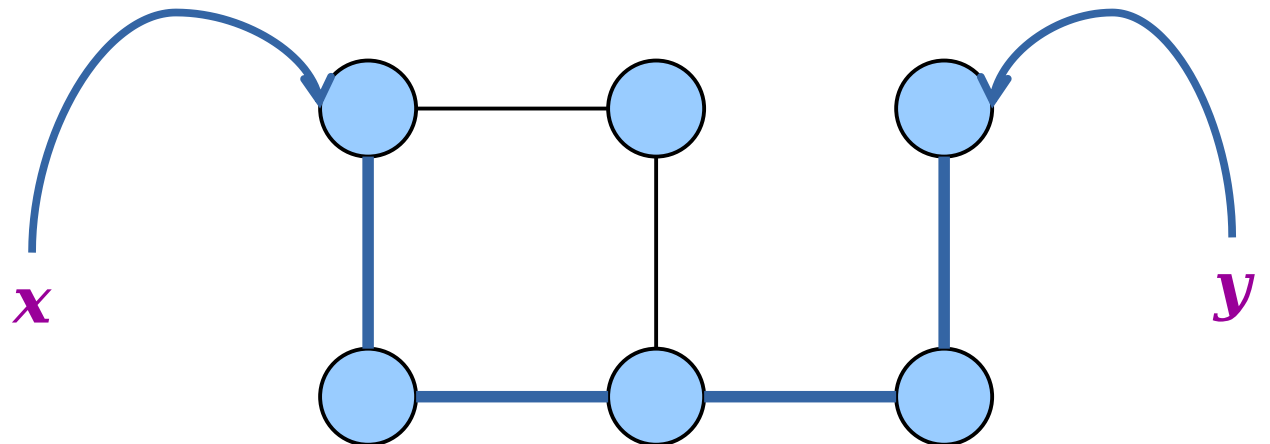
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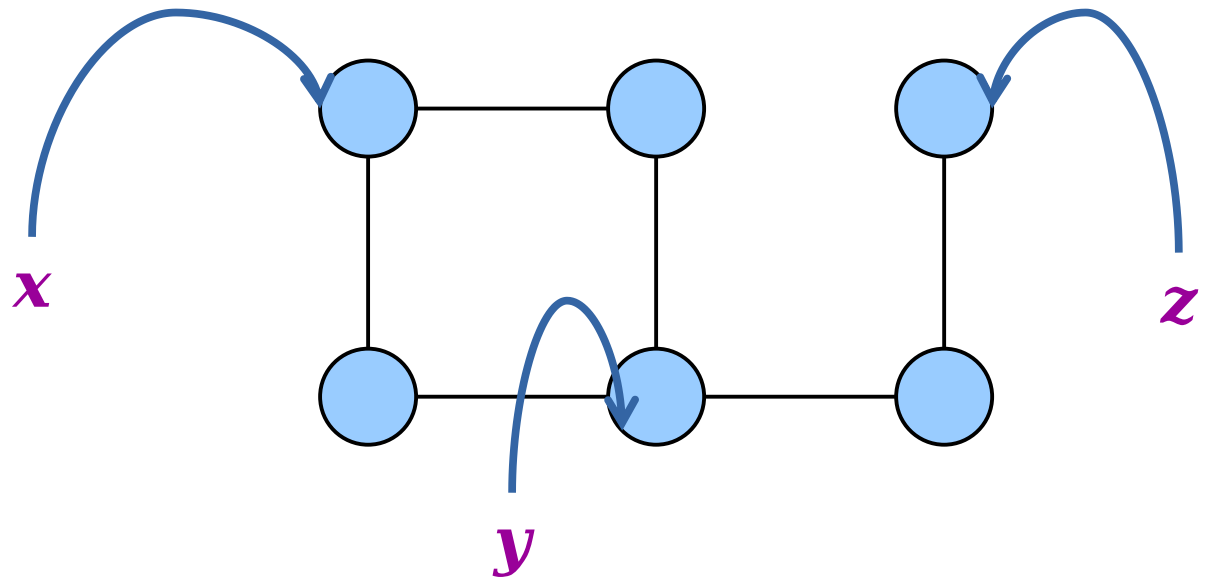
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# Connectivity

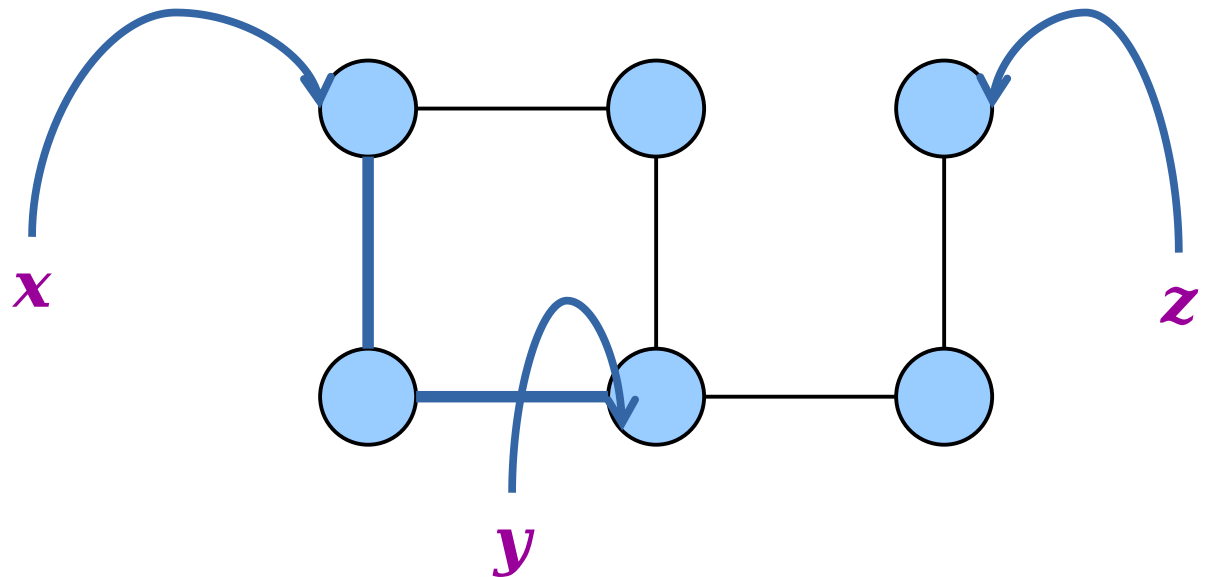
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Is it reflexive?

Is it symmetric?

Is it transitive?





**Theorem:** Let  $G = (V, E)$  be a graph. Then the connectivity relation over  $V$  is an equivalence relation.

**Proof:** Consider an arbitrary graph  $G = (V, E)$ . We will prove that the connectivity relation over  $V$  is reflexive, symmetric, and transitive.

To show that connectivity is reflexive, consider any  $v \in V$ . Then the singleton path  $v$  is a path from  $v$  to itself. Therefore,  $v$  is connected to itself, as required.

To show that connectivity is symmetric, consider any  $x, y \in V$  where  $x$  is connected to  $y$ . We need to show that  $y$  is connected to  $x$ . Since  $x$  is connected to  $y$ , there is some path  $x, v_1, \dots, v_n, y$  from  $x$  to  $y$ . Then  $y, v_n, \dots, v_1, x$  is a path from  $y$  back to  $x$ , so  $y$  is connected to  $x$ .

Finally, to show that connectivity is transitive, let  $x, y, z \in V$  be arbitrary nodes where  $x$  is connected to  $y$  and  $y$  is connected to  $z$ . We will prove that  $x$  is connected to  $z$ . Since  $x$  is connected to  $y$ , there is a path  $x, u_1, \dots, u_n, y$  from  $x$  to  $y$ . Since  $y$  is connected to  $z$ , there is a path  $y, v_1, \dots, v_k, z$  from  $y$  to  $z$ . Then the path  $x, u_1, \dots, u_n, y, v_1, \dots, v_k, z$  goes from  $x$  to  $z$ . Thus  $x$  is connected to  $z$ , as required. ■

# Putting Things Together

Earlier, we defined the connected component of a node  $v$  to be

$$[v] = \{ x \in V \mid v \text{ is connected to } x \}$$

Connectivity is an equivalence relation! So what's the equivalence class of a node  $v$  with respect to connectivity?

$$[v]_{\text{conn}} = \{ x \in V \mid v \text{ is connected to } x \}$$

***Connected components are equivalence classes of the connectivity relation!***

**Theorem:** If  $G = (V, E)$  is a graph, then every node in  $G$  belongs to exactly one connected component of  $G$ .

**Proof:** Let  $G = (V, E)$  be an arbitrary graph and let  $v \in V$  be any node in  $G$ . The connected components of  $G$  are just the equivalence classes of the connectivity relation in  $G$ . The Fundamental Theorem of Equivalence Relations guarantees that  $v$  belongs to exactly one equivalence class of the connectivity relation. Therefore,  $v$  belongs to exactly one connected component in  $G$ . ■

**Time out for announcements!**

# Problem Set 3

- Due tomorrow (Thursday) at 11:59pm PDT.
- Use a late period to extend this to Saturday at 11:59pm PDT.
- Any last questions? Come to office hours or ask on Campuswire.

# Midterm

- Thursday July 23<sup>rd</sup>
- Will cover material up to Monday's lecture (psets 1, 2, 3).
- 24-hour window to start the exam. Begins at 9:30AM PDT on Thursday July 23<sup>rd</sup>.
- Once you click start on Gradescope, Gradescope will give you access to the exam. You'll have 3 hours to complete the exam, plus 15 minutes to upload your exam to Gradescope.
- Please make sure any OAE letters get sent to the staff mailing list as soon as possible.