Eric Roberts CS 106B

Handout #45 February 27, 2013

Graph Algorithms

	Outline
Graph Algorithms	 A review the graphtypes.h and graph.h interfaces A tour of the Pathfinder assignment Examples of depth-first and breadth-first search Dijkstra's shortest-path algorithm Kruskal's minimum-spanning-tree algorithm
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- One of the most useful algorithms for computing the shortest paths in a graph was developed by Edsgar W. Dijkstra in 1959.
- The strategy is similar to the breadth-first search algorithm you used to implement the word-ladder program in Assignment #2. The major difference are:
- The queue used to hold the paths delivers items in increasing order of total cost rather than in the traditional first-in/first-out order. Such queues are called *priority queues*.
- The algorithm keeps track of all nodes to which the total distance has already been fixed. Distances are fixed whenever you dequeue a path from the priority queue.





Kruskal's Algorithm

- In many cases, finding the shortest path is not as important as as minimizing the cost of a network as a whole. A set of arcs that connects every node in a graph at the smallest possible cost is called a *minimum spanning tree*.
- The following algorithm for finding a minimum spanning tree was developed by Joseph Kruskal in 1956:
 - Start with a new empty graph with the same nodes as the original one but an empty set of arcs.
 - Sort all the arcs in the graph in order of increasing cost.
 - Go through the arcs in order and add each one to the new graph if the endpoints of that arc are not already connected by a path.
- This process can be made more efficient by maintaining sets of nodes in the new graph, as described on the next slide.

Combining Sets in Kruskal's Algorithm

- Implementing the Pathfinder version of Kruskal's algorithm requires you need to build a new graph containing the spanning tree. As you do, you will generate sets of disconnected graphs.
- When you choose a new arc, there are four possibilities for the sets formed by the nodes at the endpoints:
- 1. *Neither node is yet in a set.* In this case, create a new set and add both nodes to it.
- 2. One node is in a set and the other isn't. In this case, add the new node to the same set.
- The endpoints are in different existing sets. In this case, you need to merge the two sets to create a new one containing the union of the existing ones.
- 4. *The endpoints are in the same set.* In this case, there is already a path between these two nodes, so you don't need this arc.

Exercise: Minimum Spanning Tree

Apply Kruskal's algorithm to find a minimum spanning tree:



An Application of Kruskal's Algorithm

• Suppose that you have a graph that looks like this:



• What would happen if you applied Kruskal's algorithm for finding a minimum spanning tree, assuming that you choose the arcs in a random order?