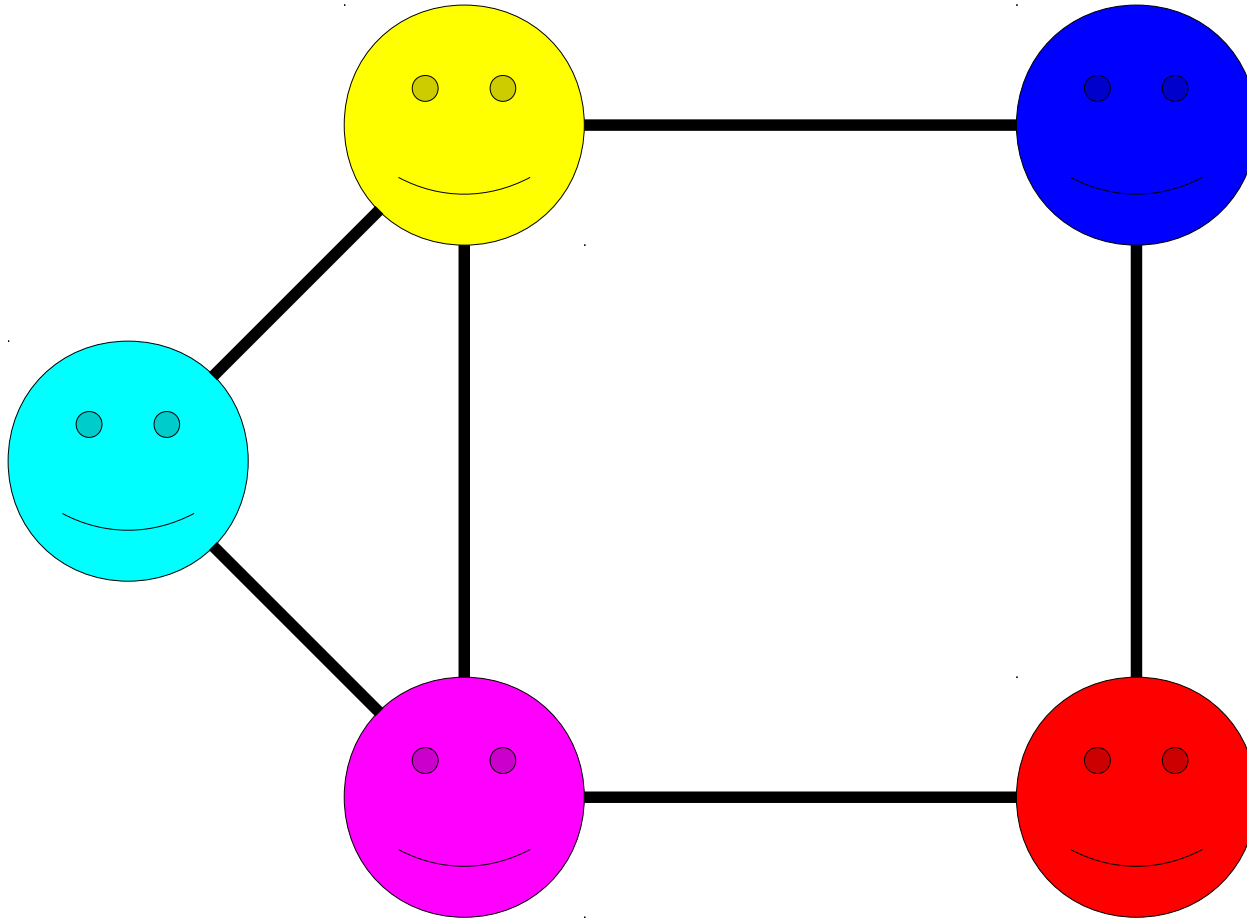
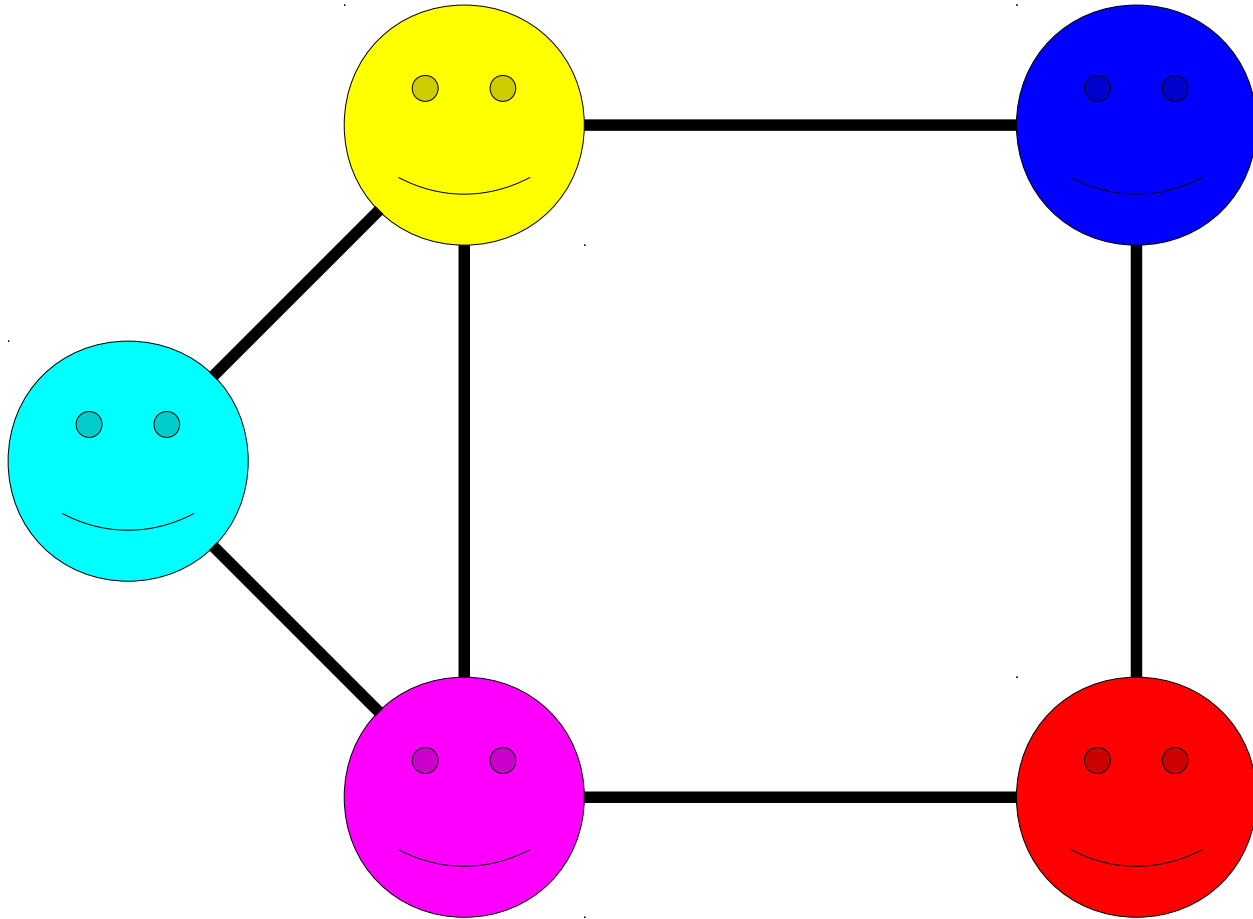


Shortest Paths

A **graph** is a mathematical structure for representing relationships.

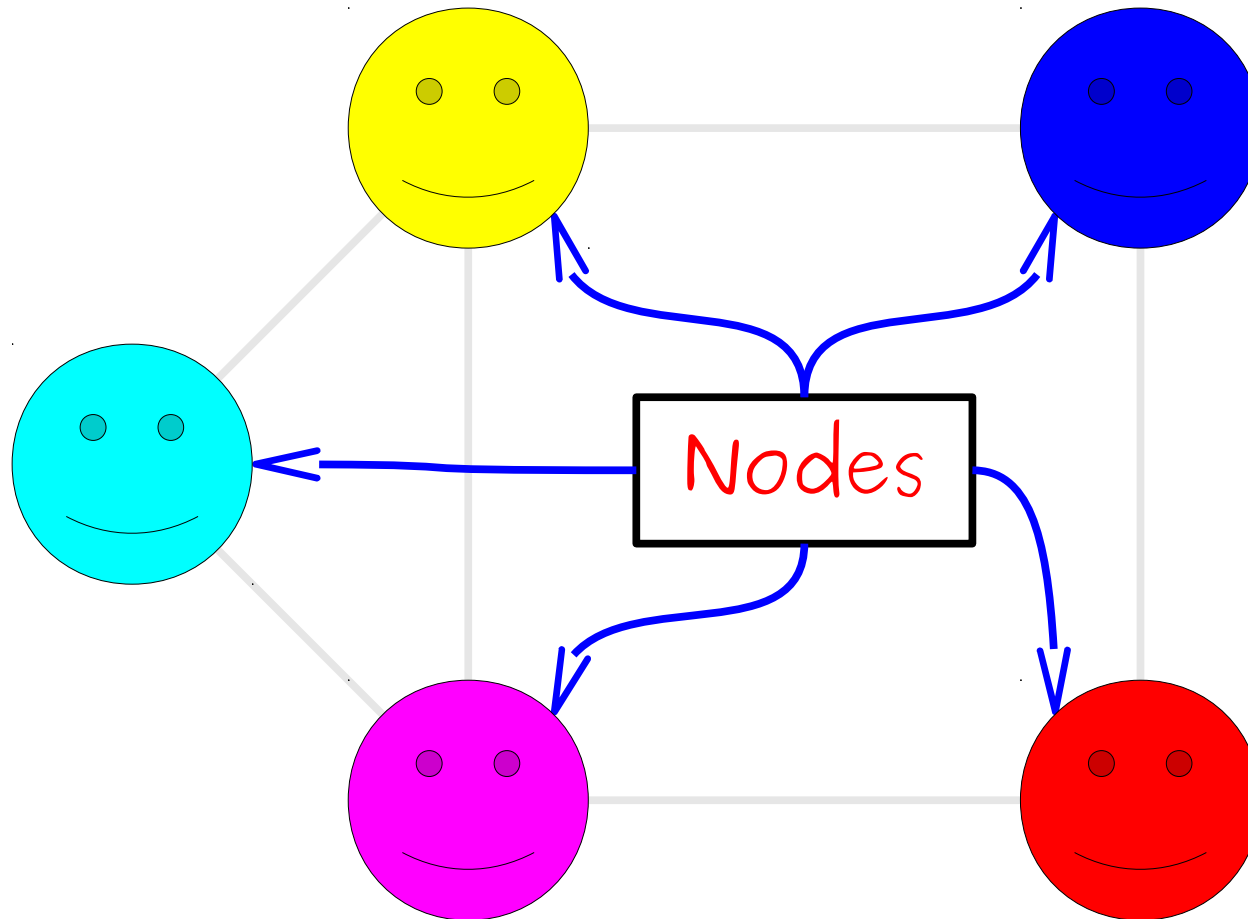


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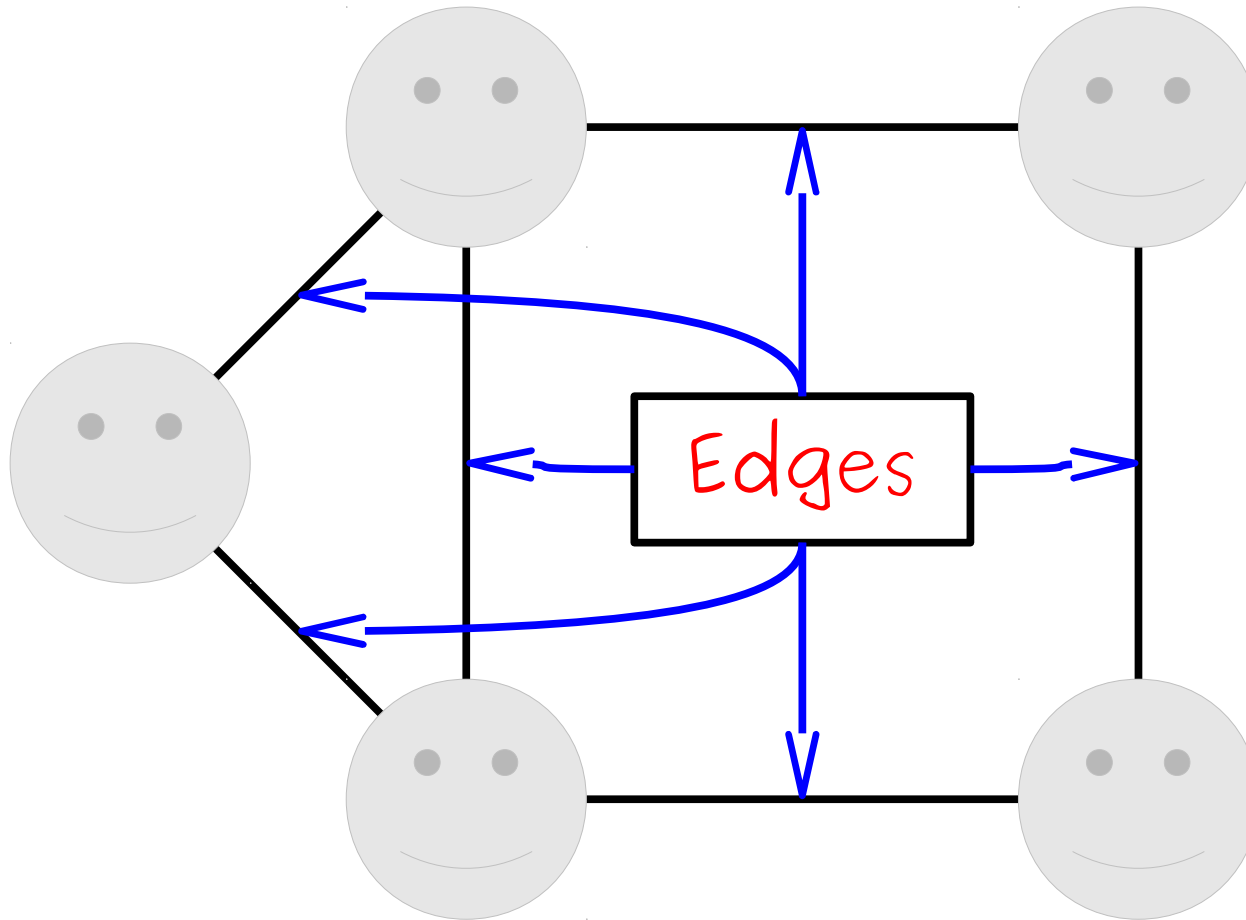
A graph consists of a set of **nodes** connected by **edges**.

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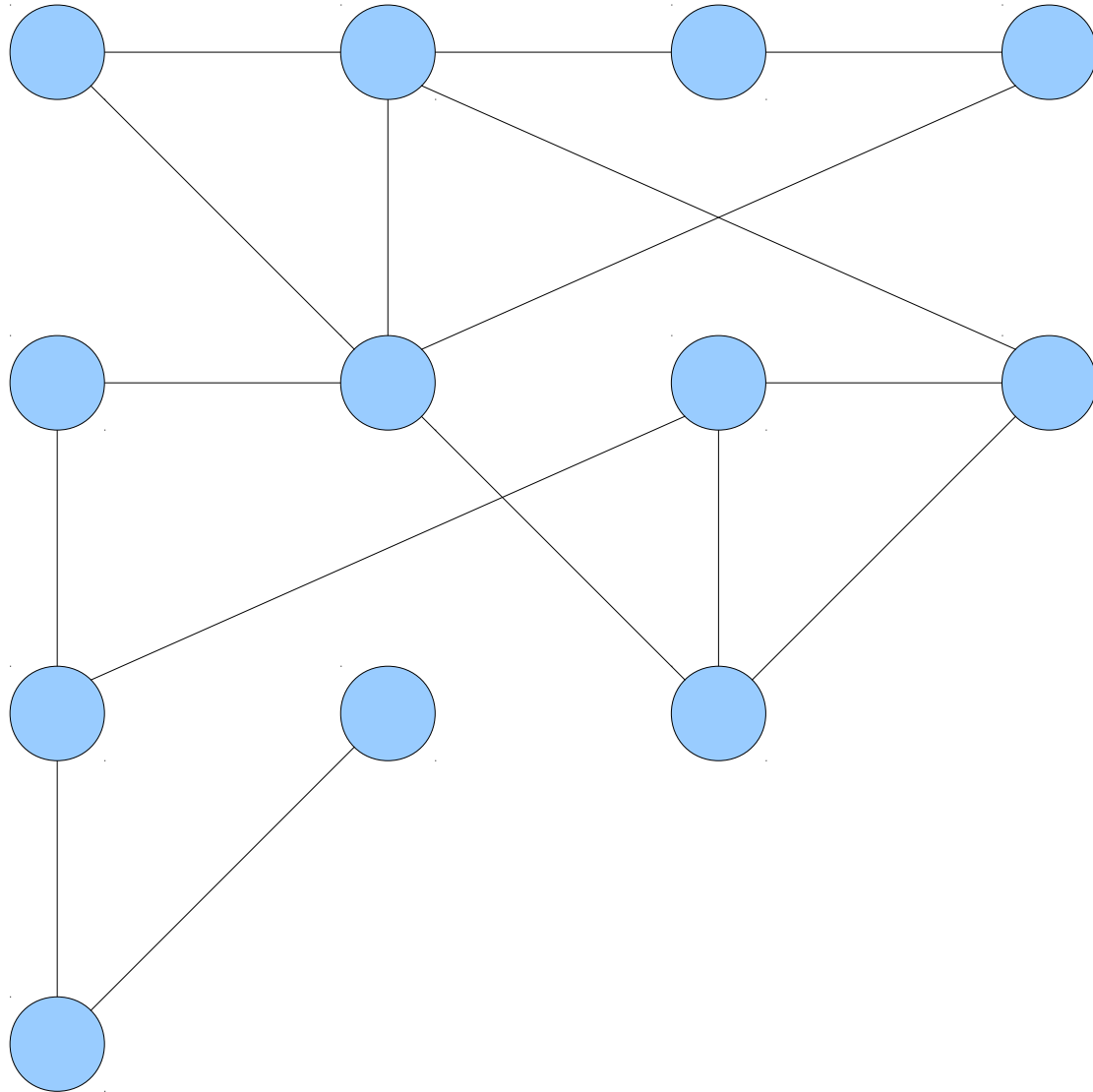
A **graph** is a mathematical structure for representing relationships.



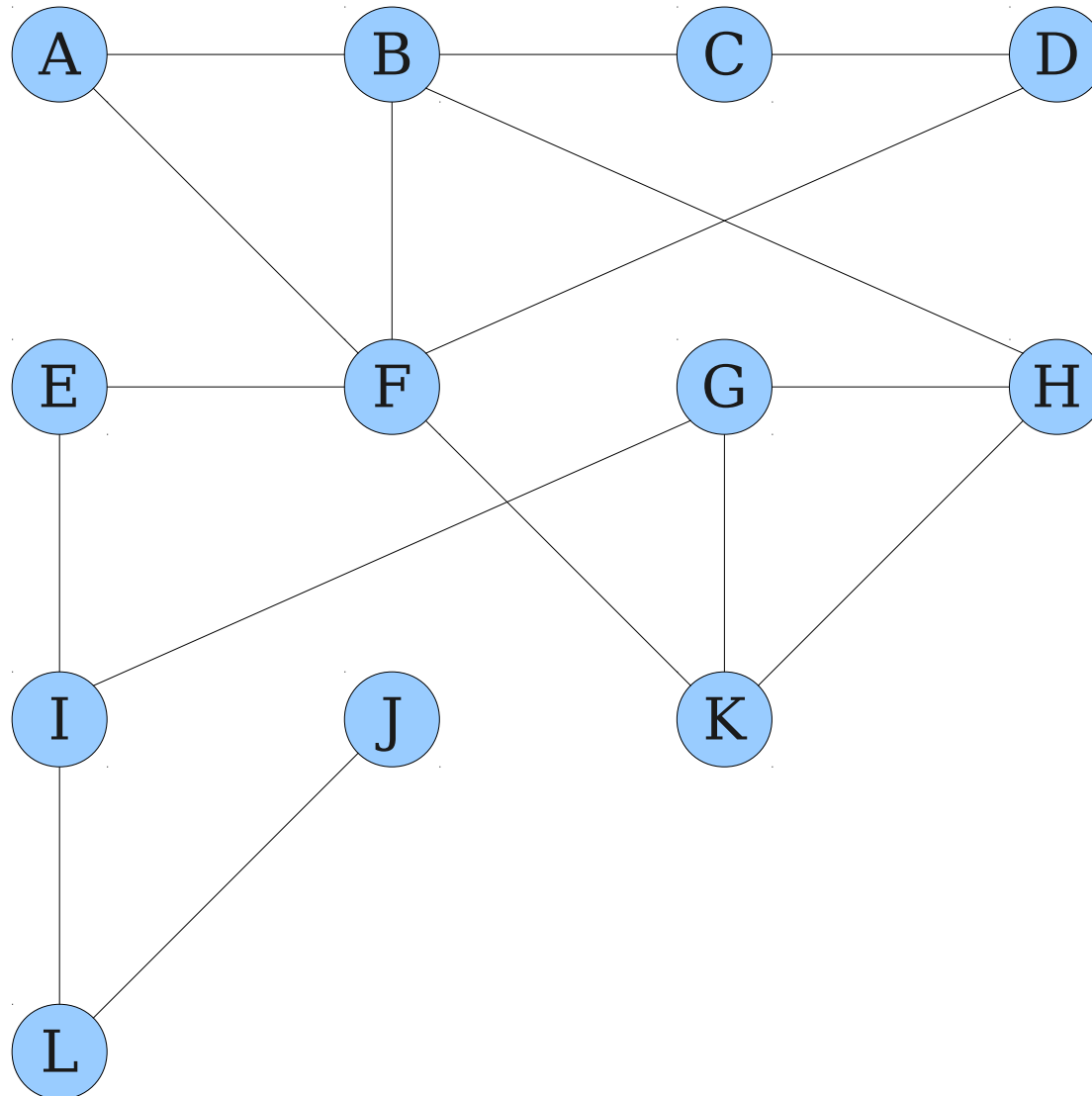
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Shortest Paths

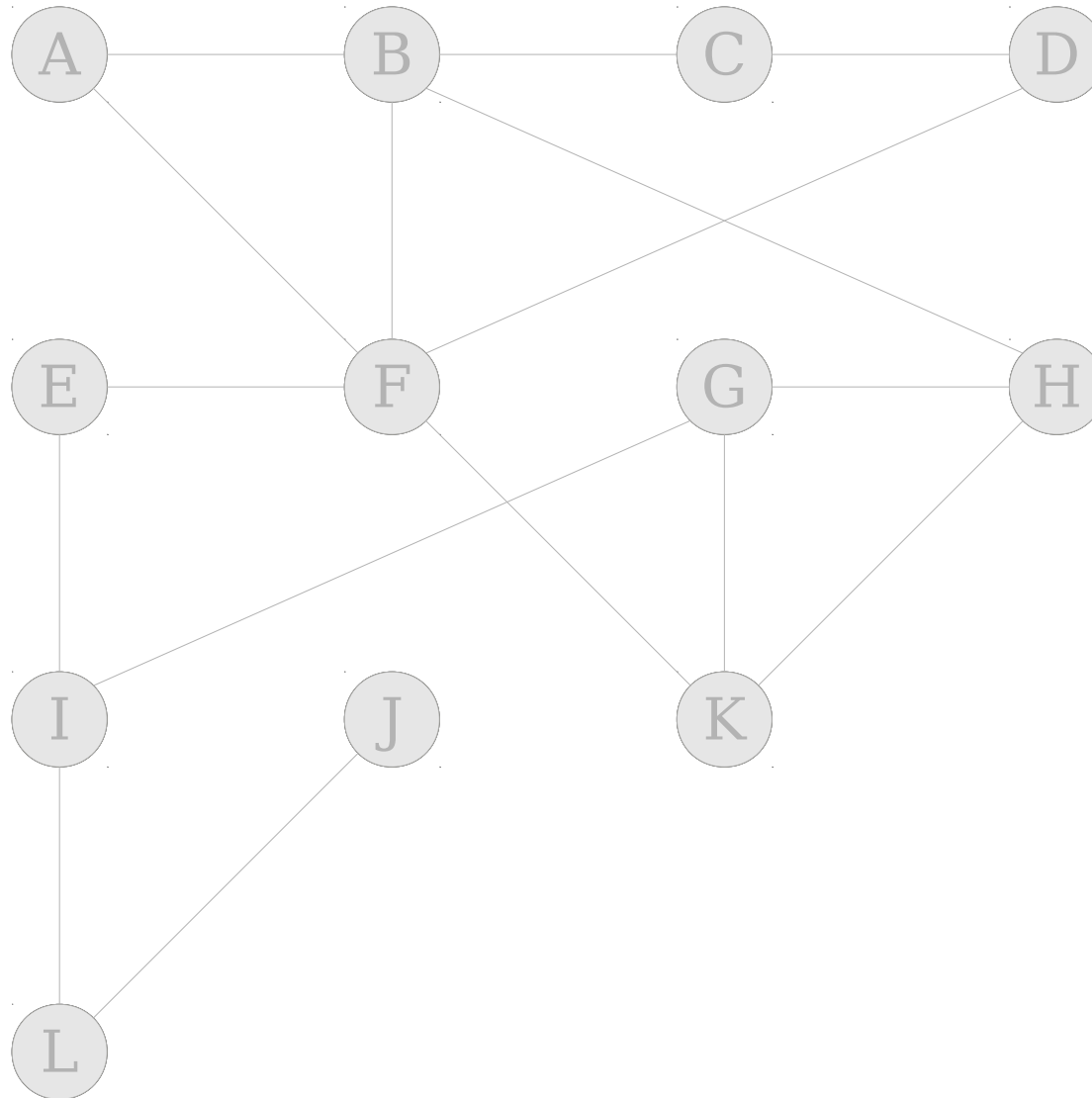
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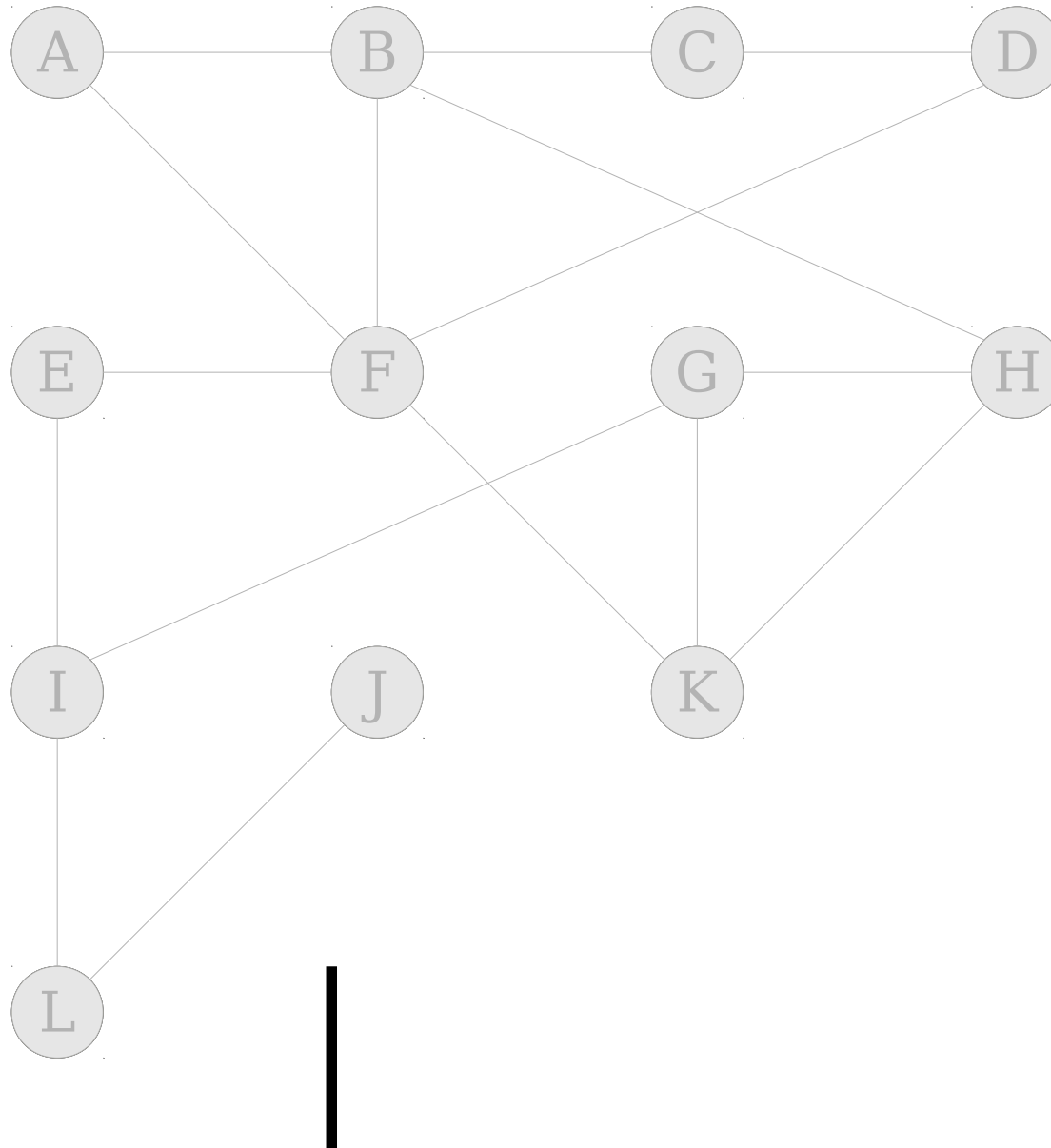
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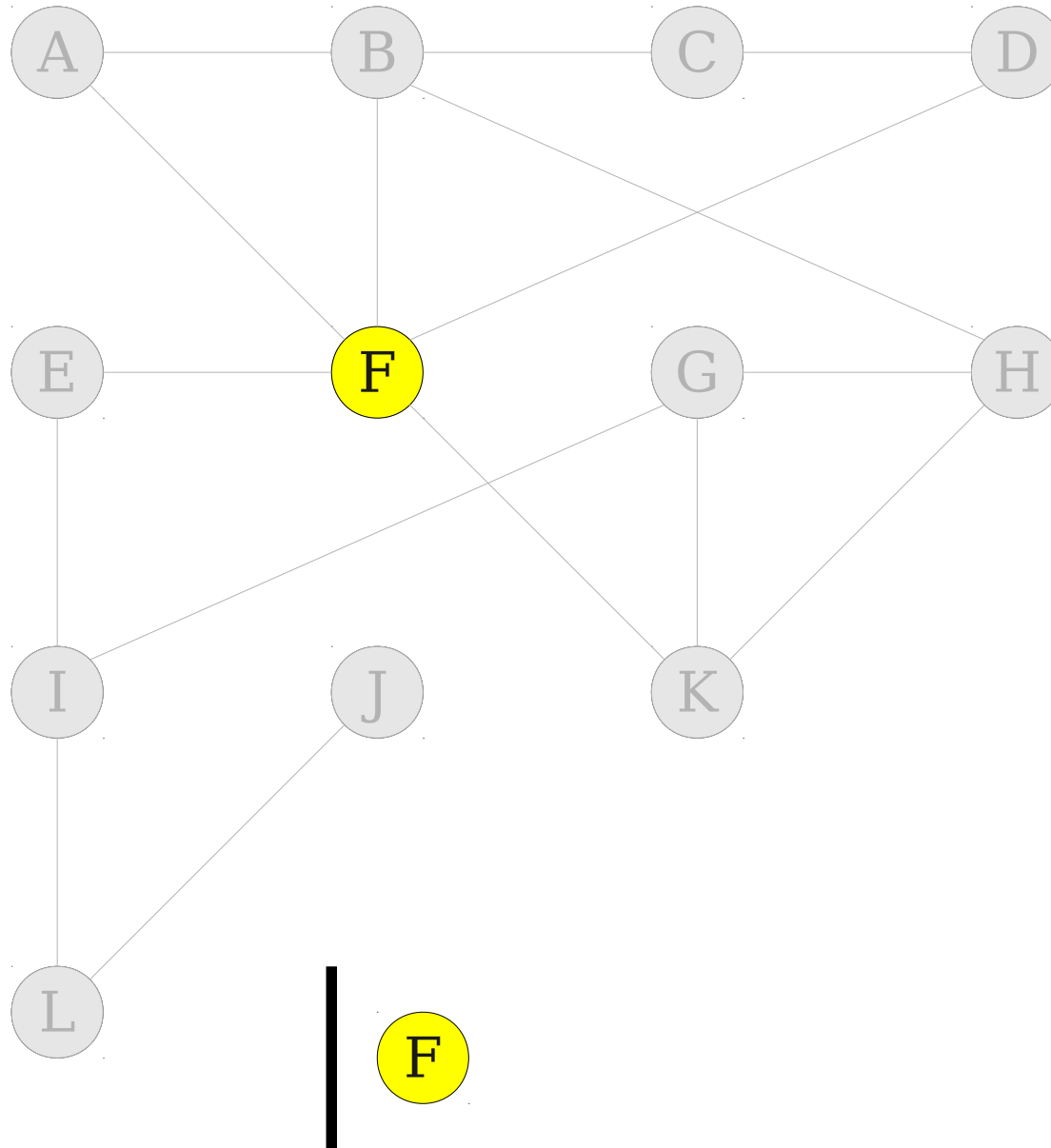
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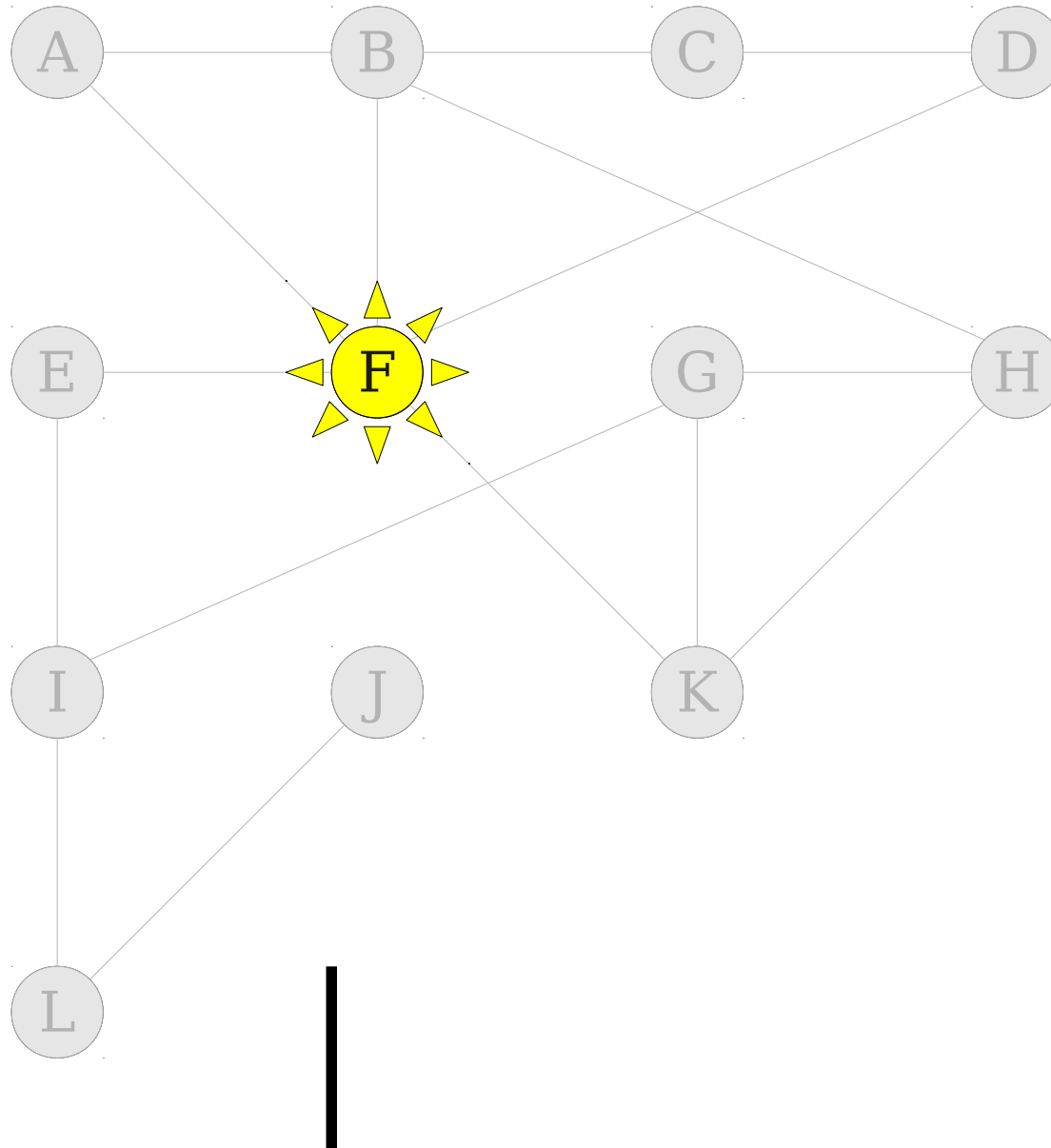
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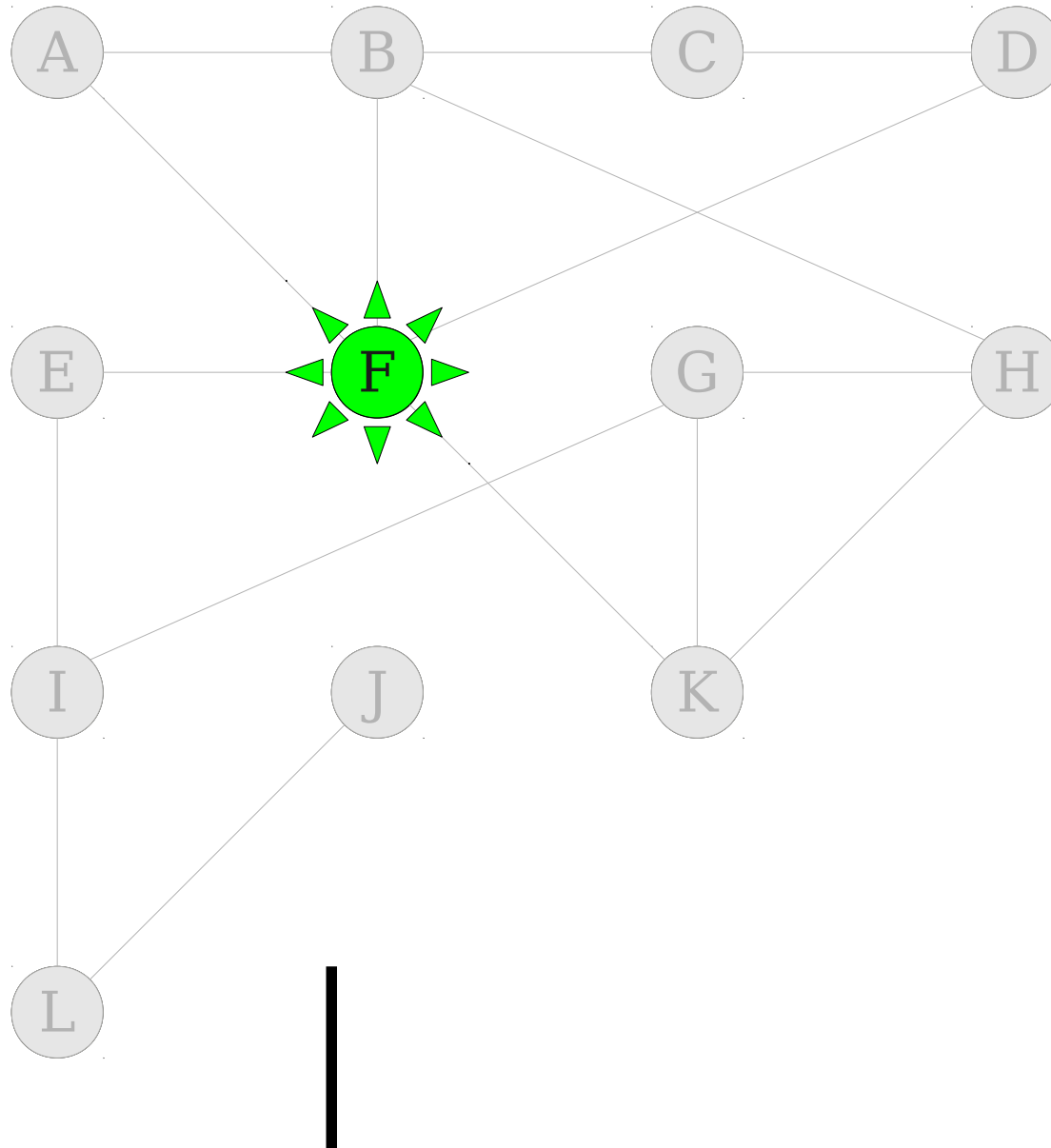
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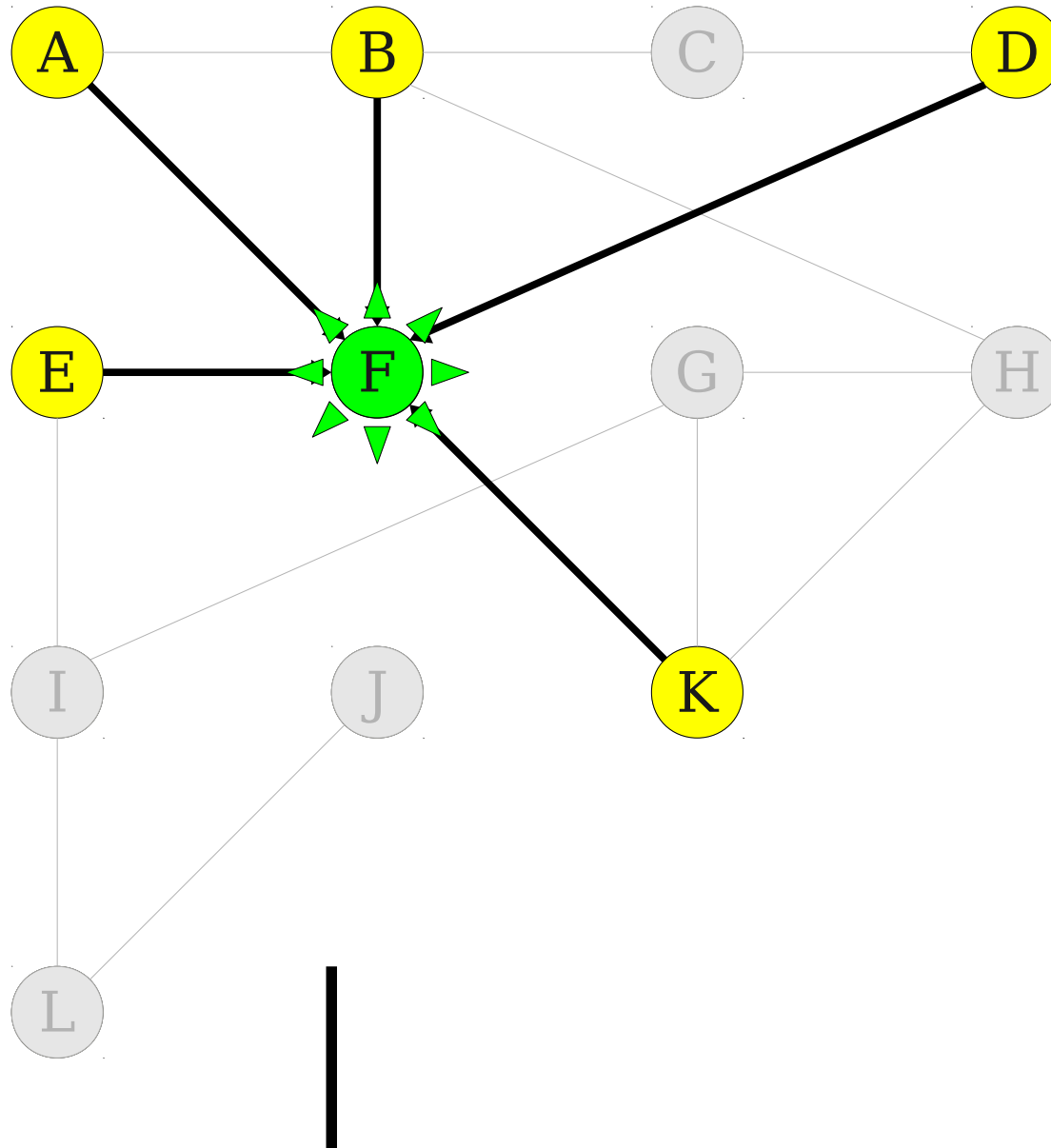
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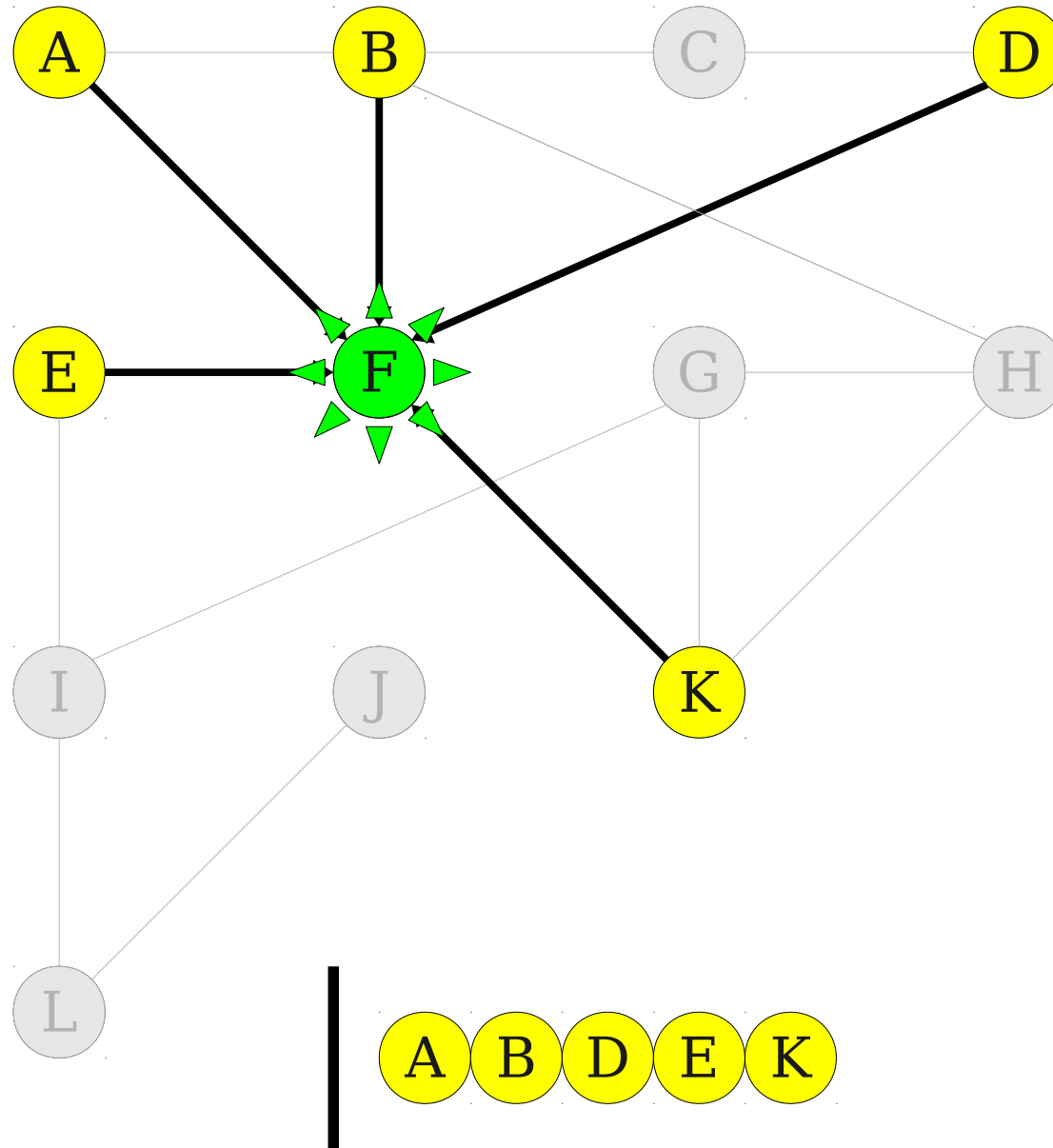
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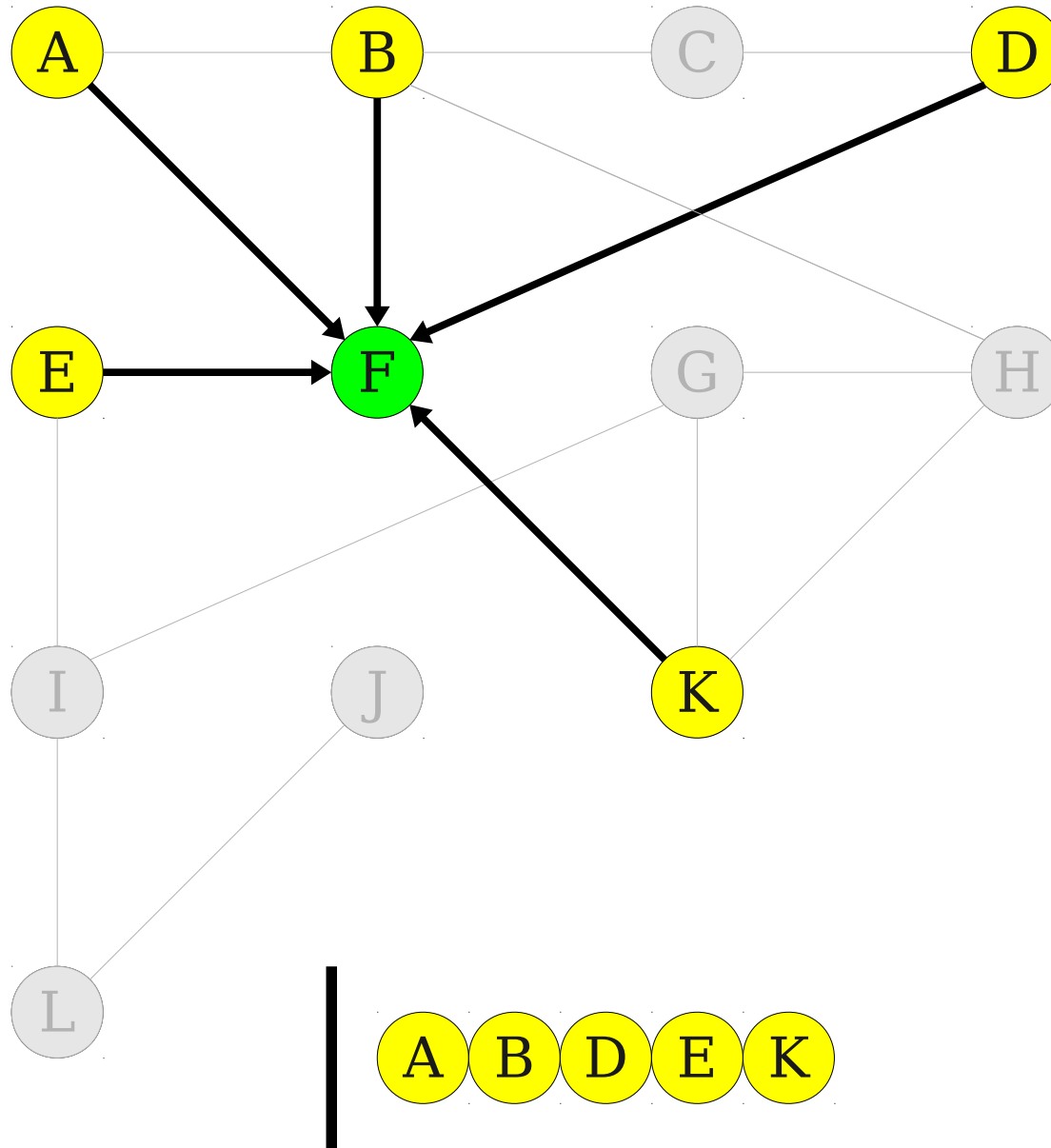
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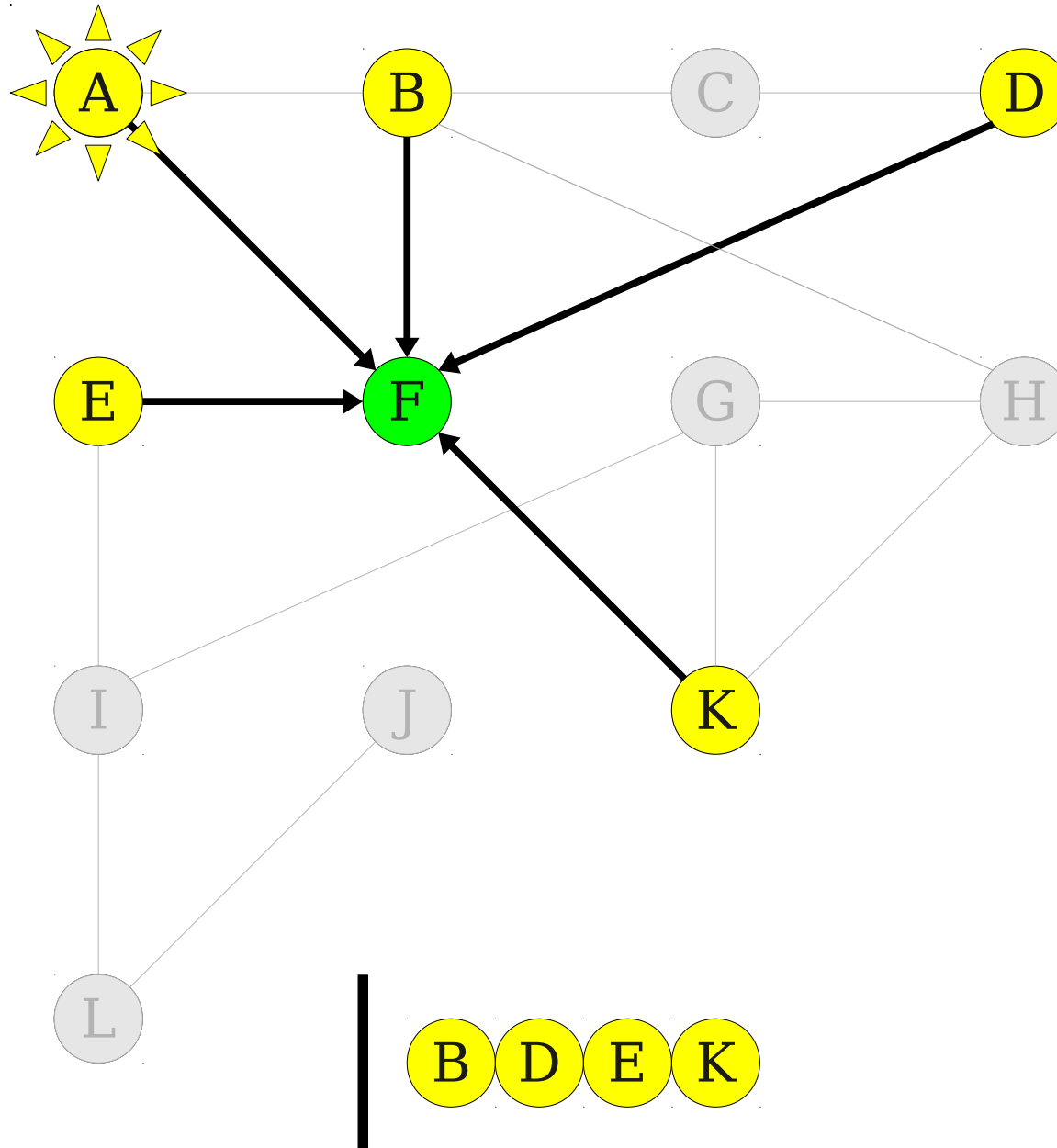
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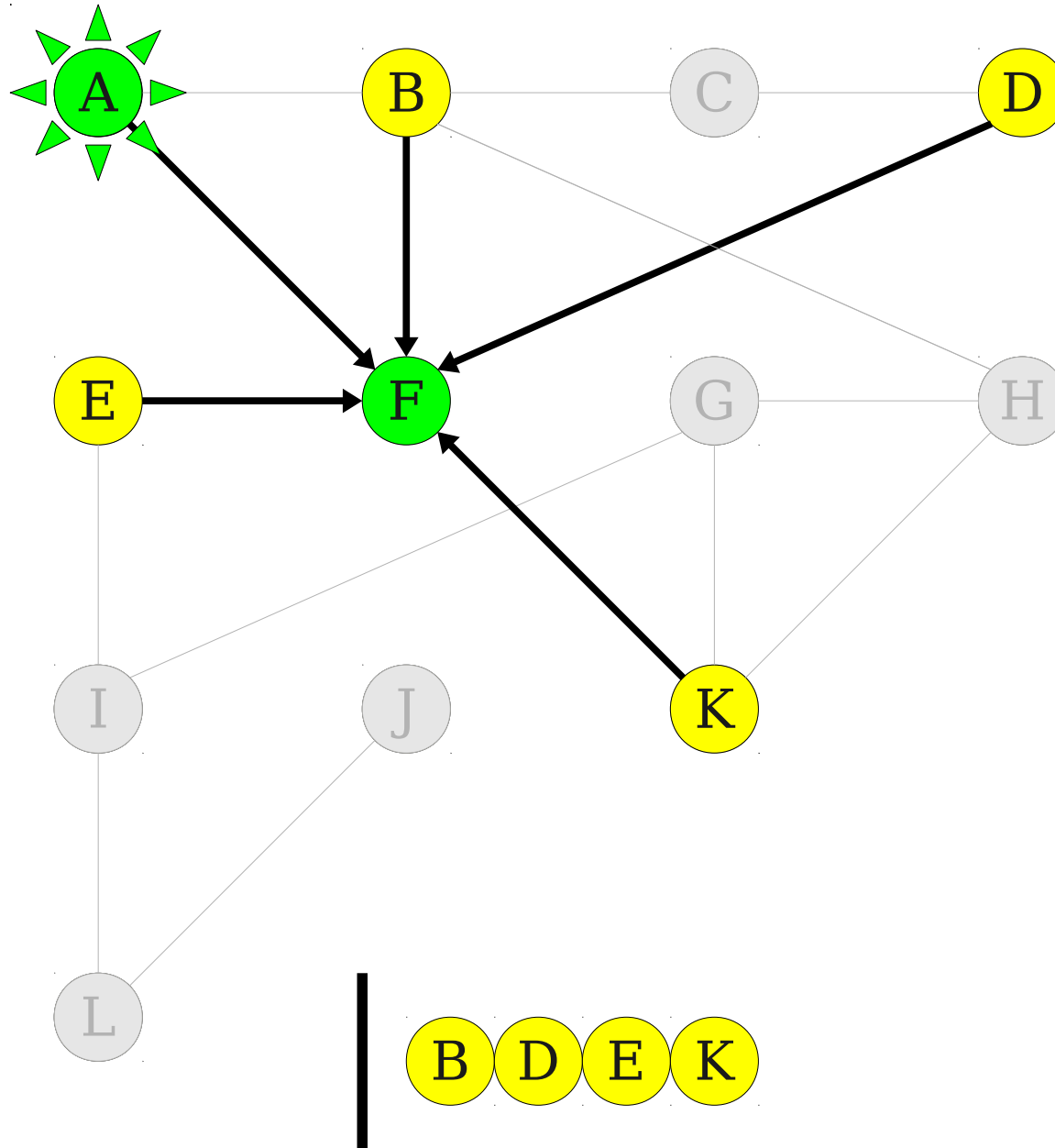
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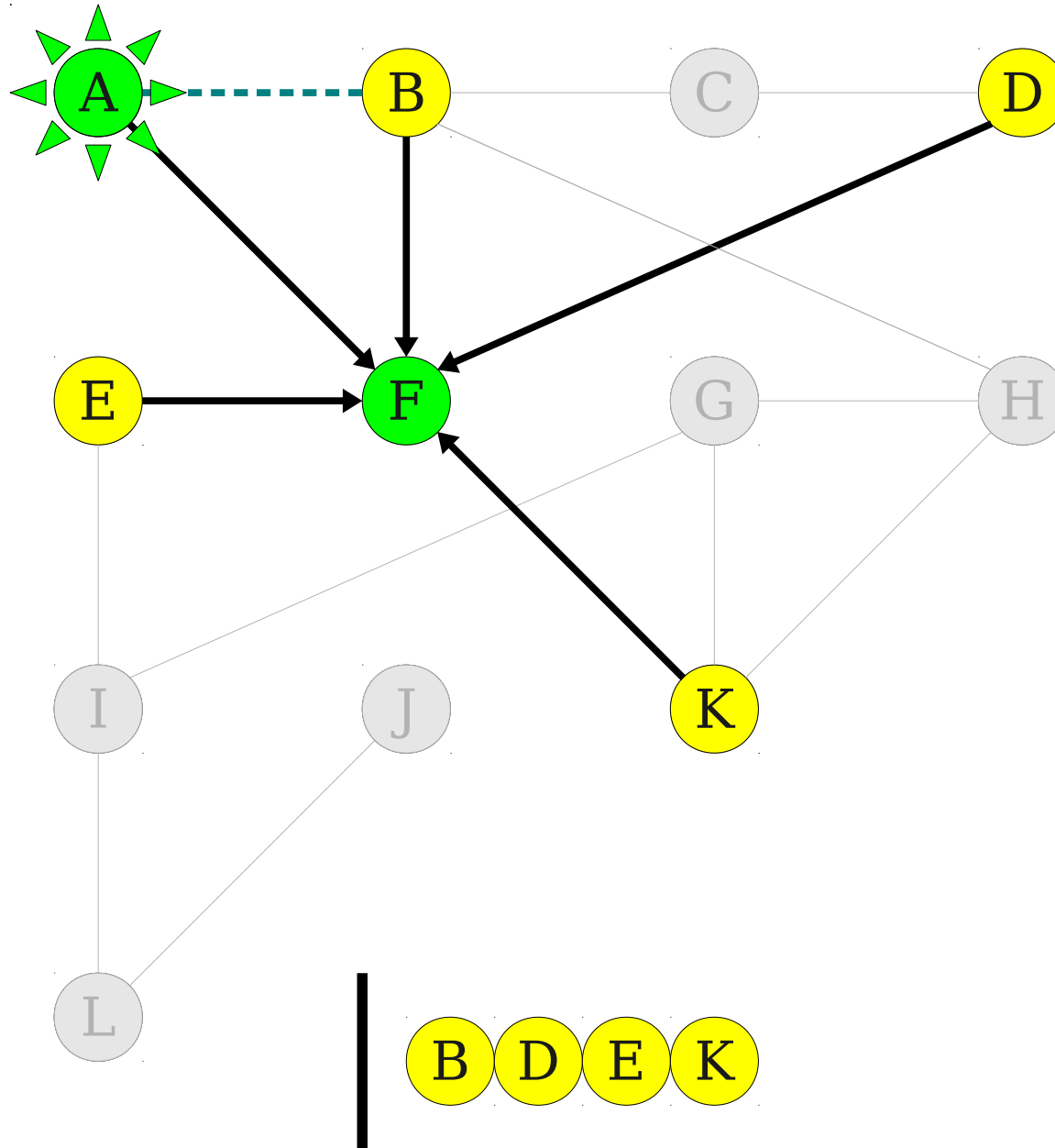
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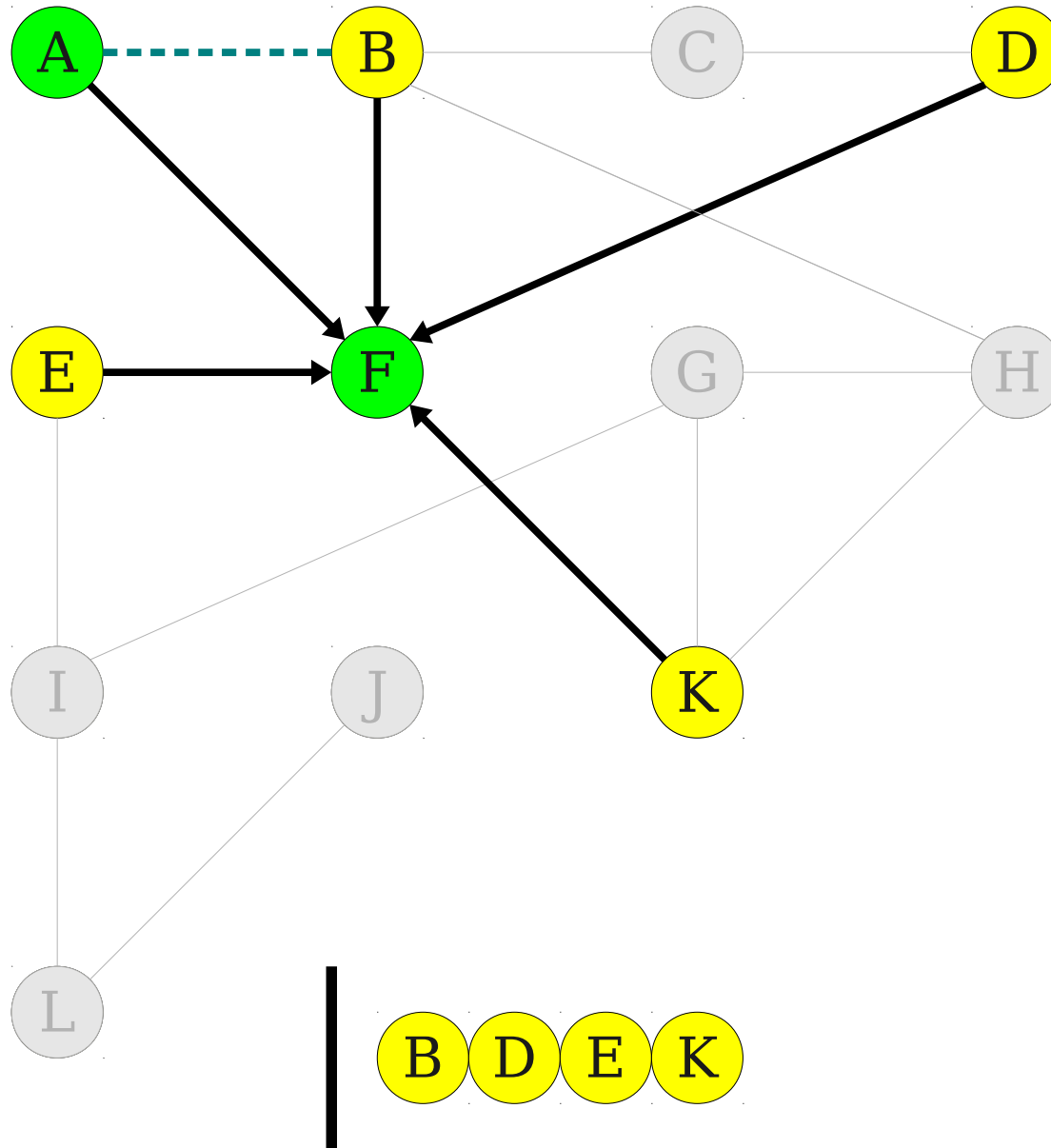
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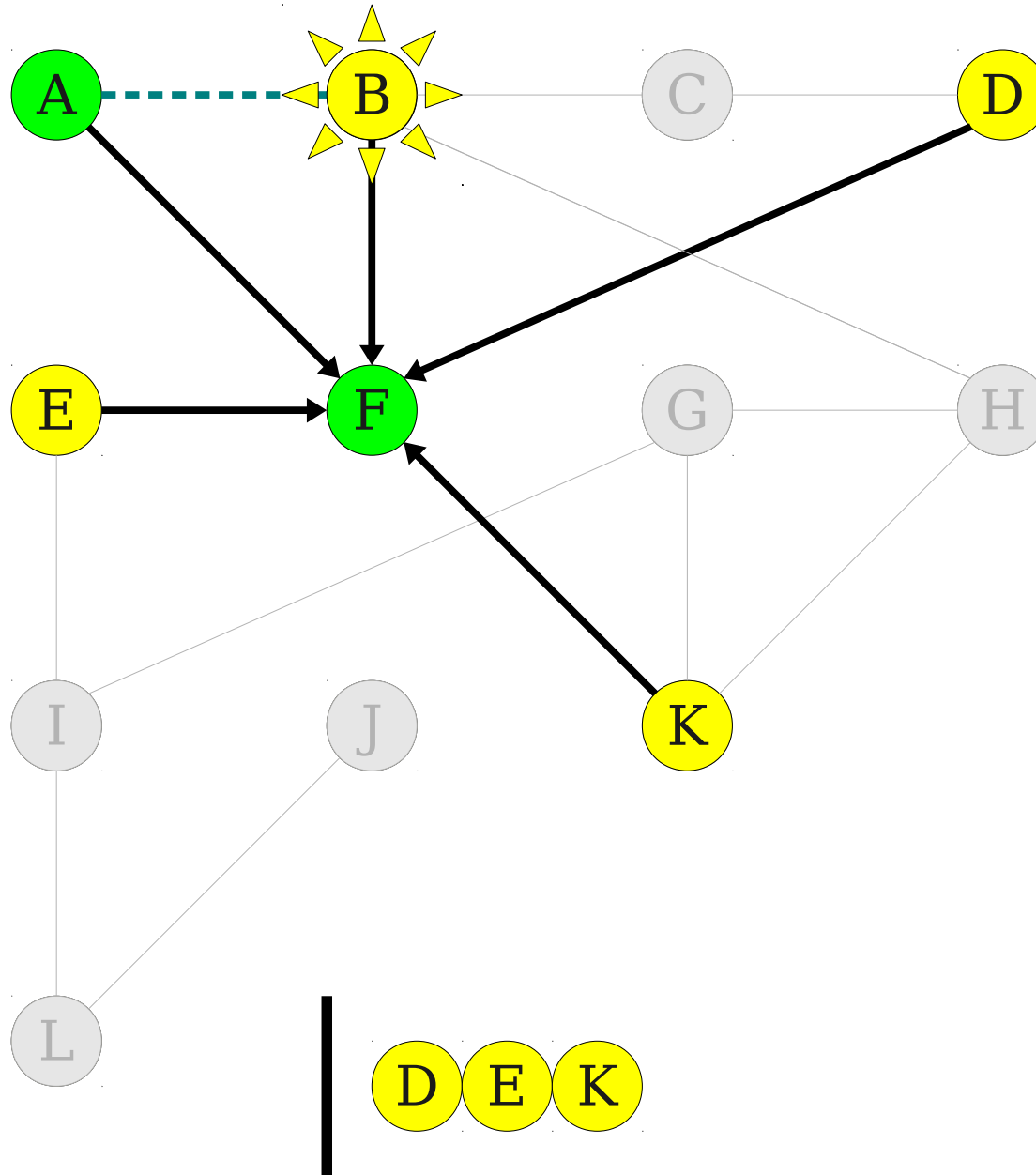
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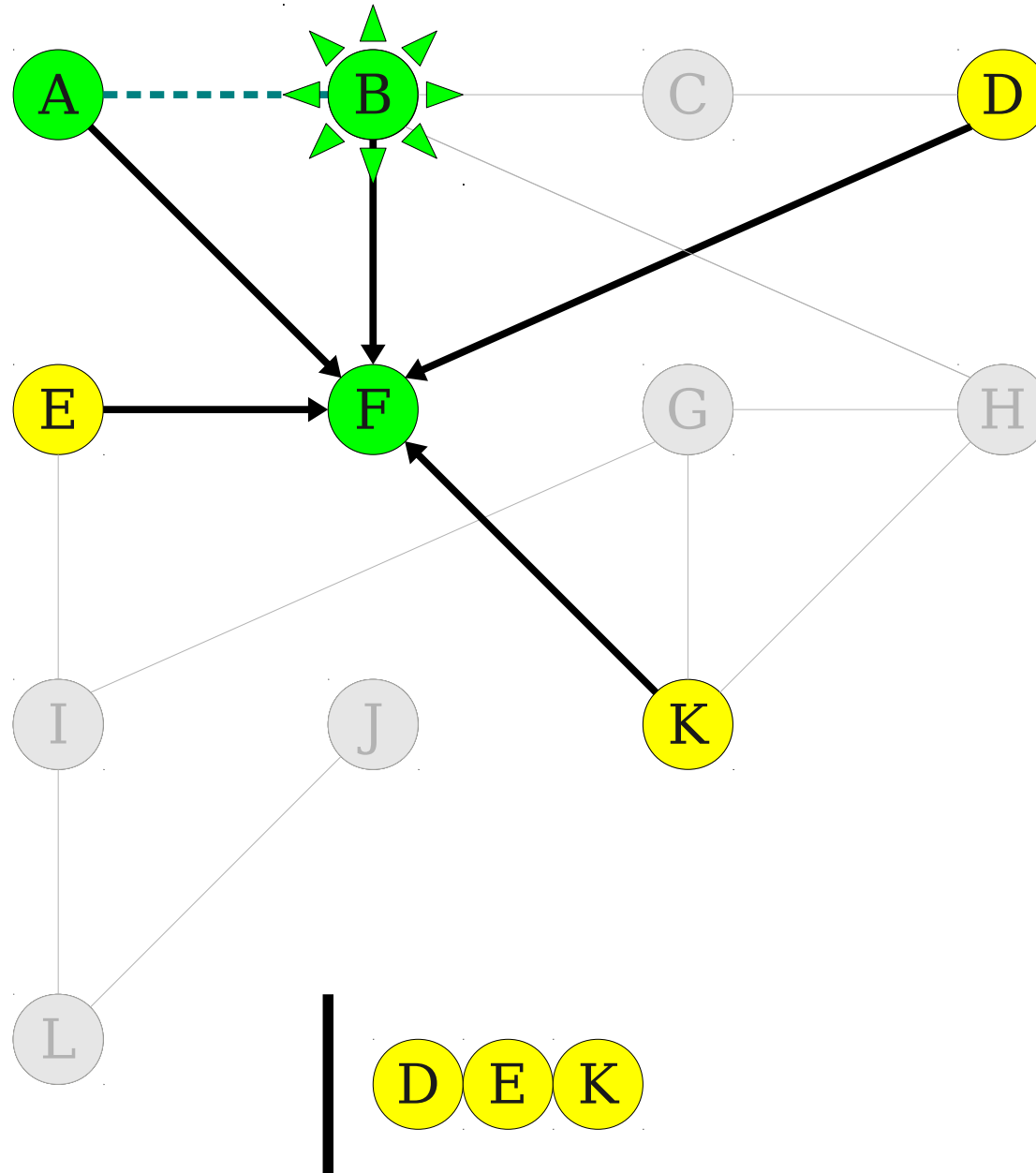
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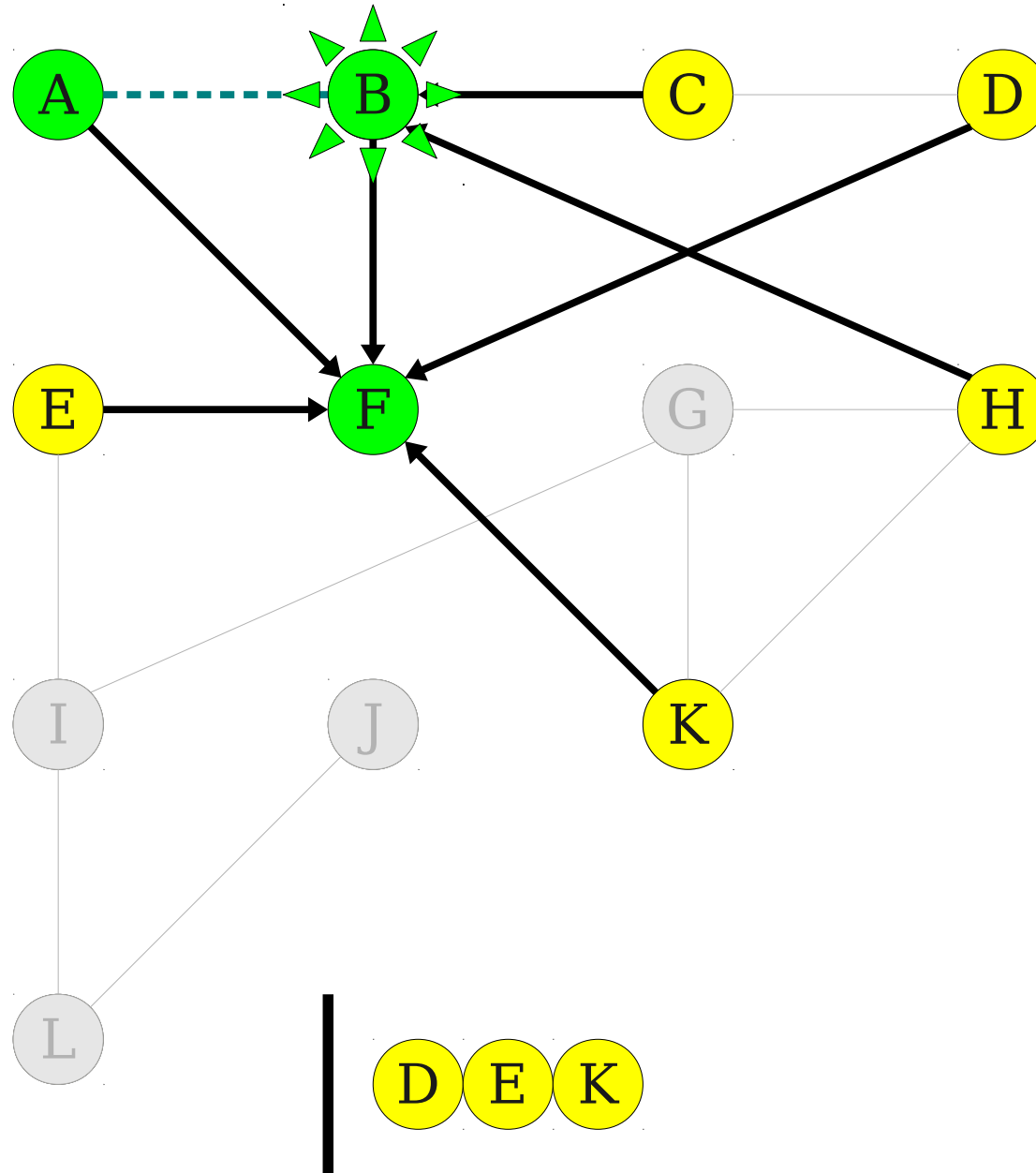
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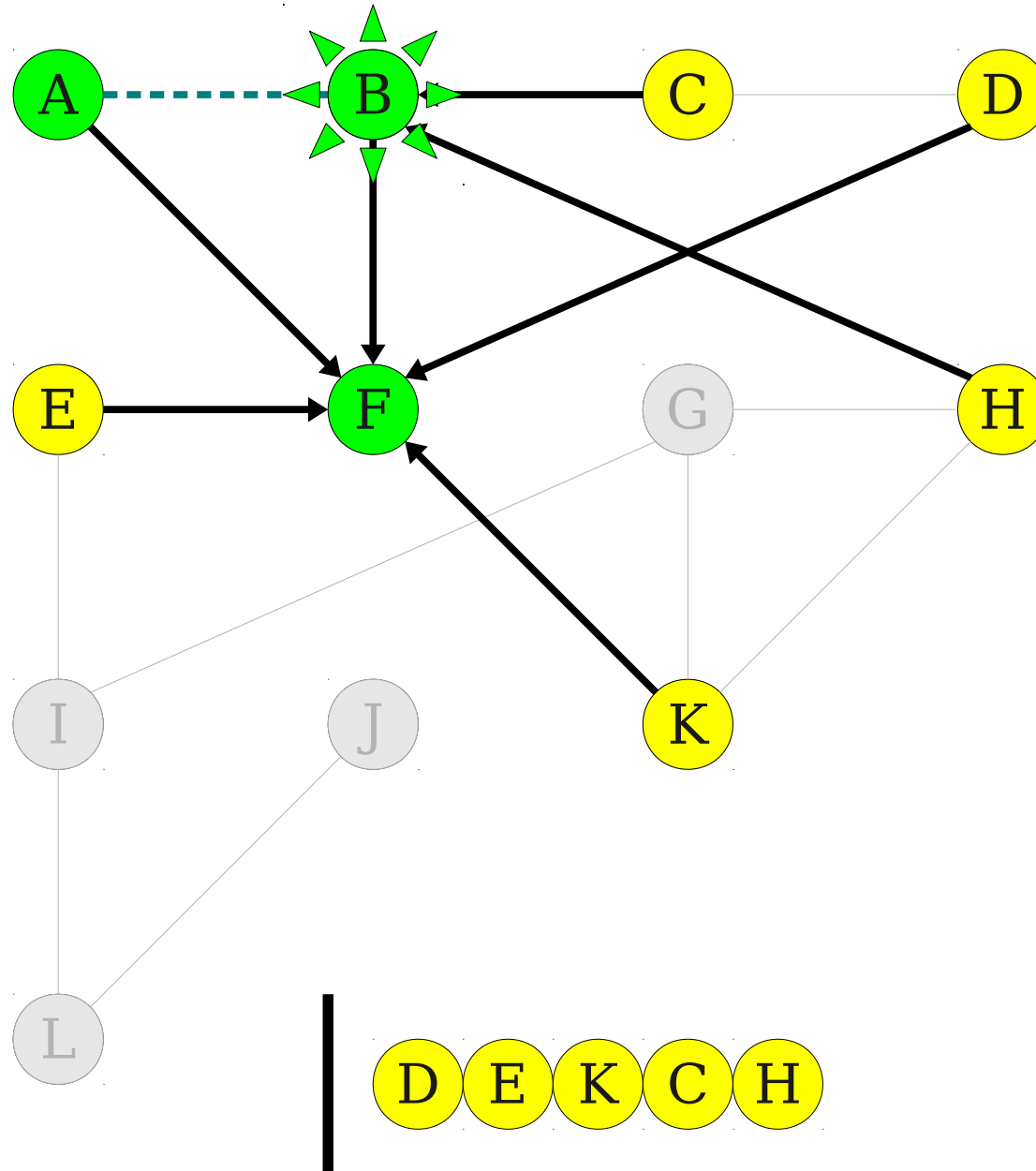
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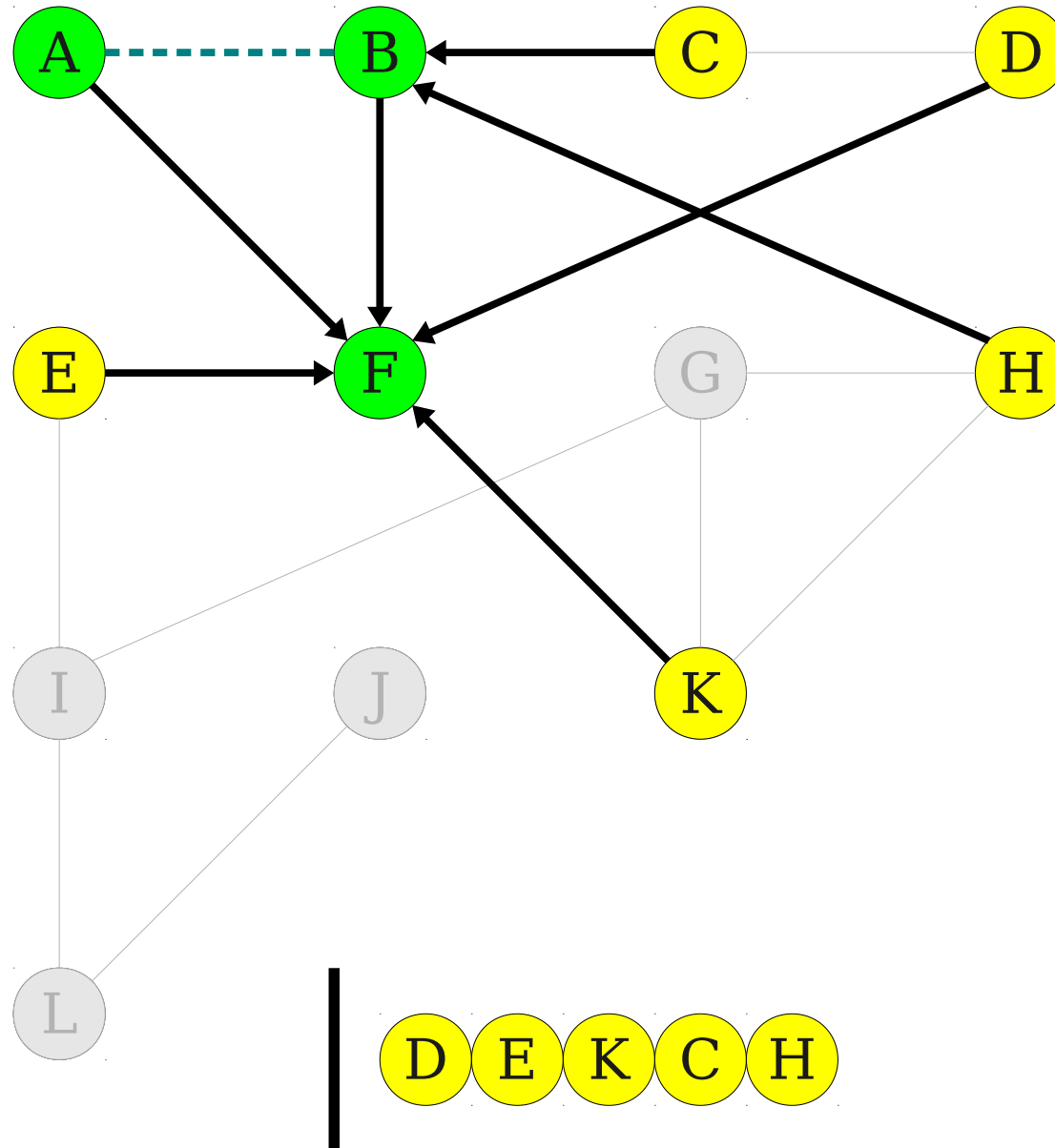
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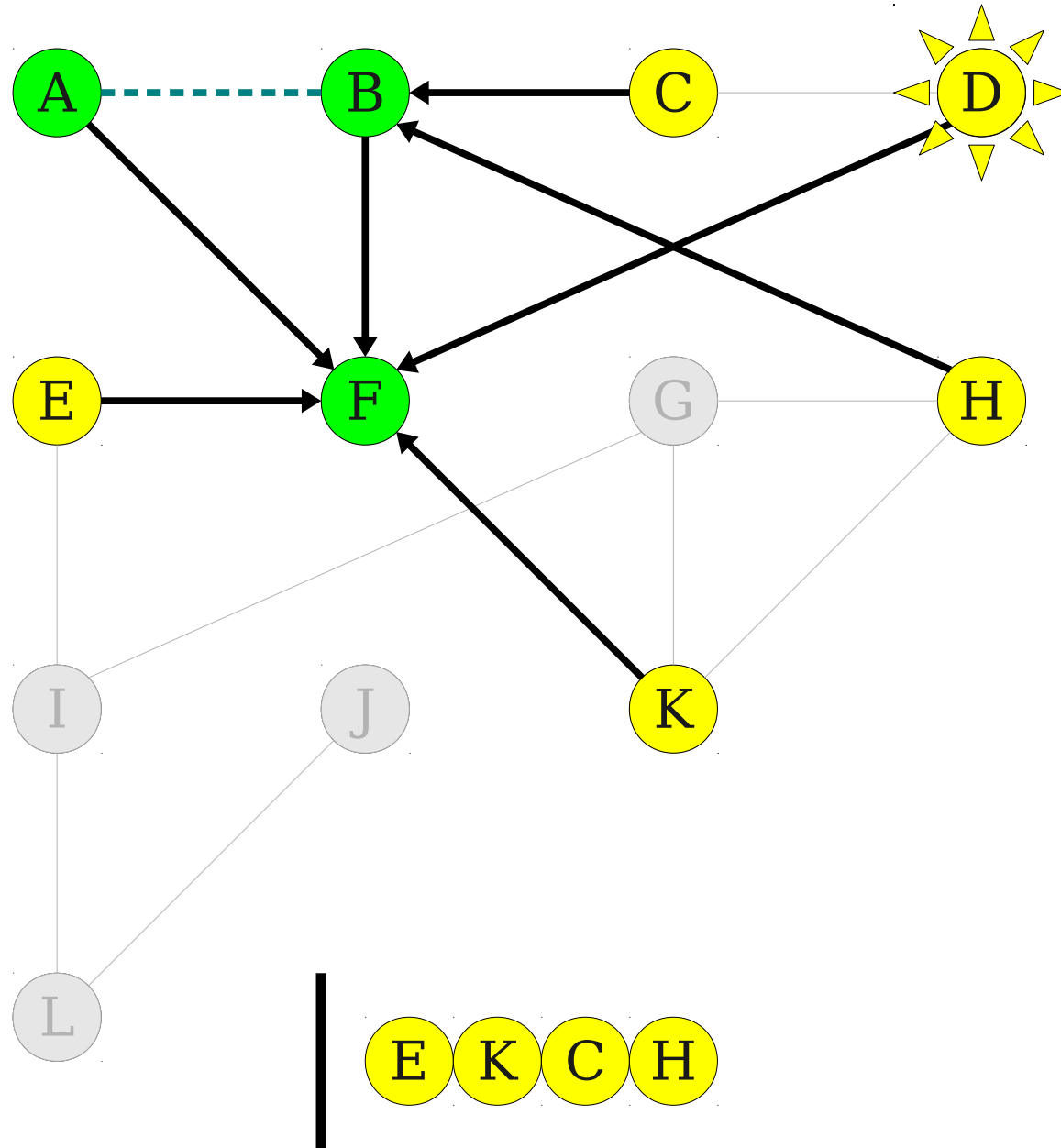
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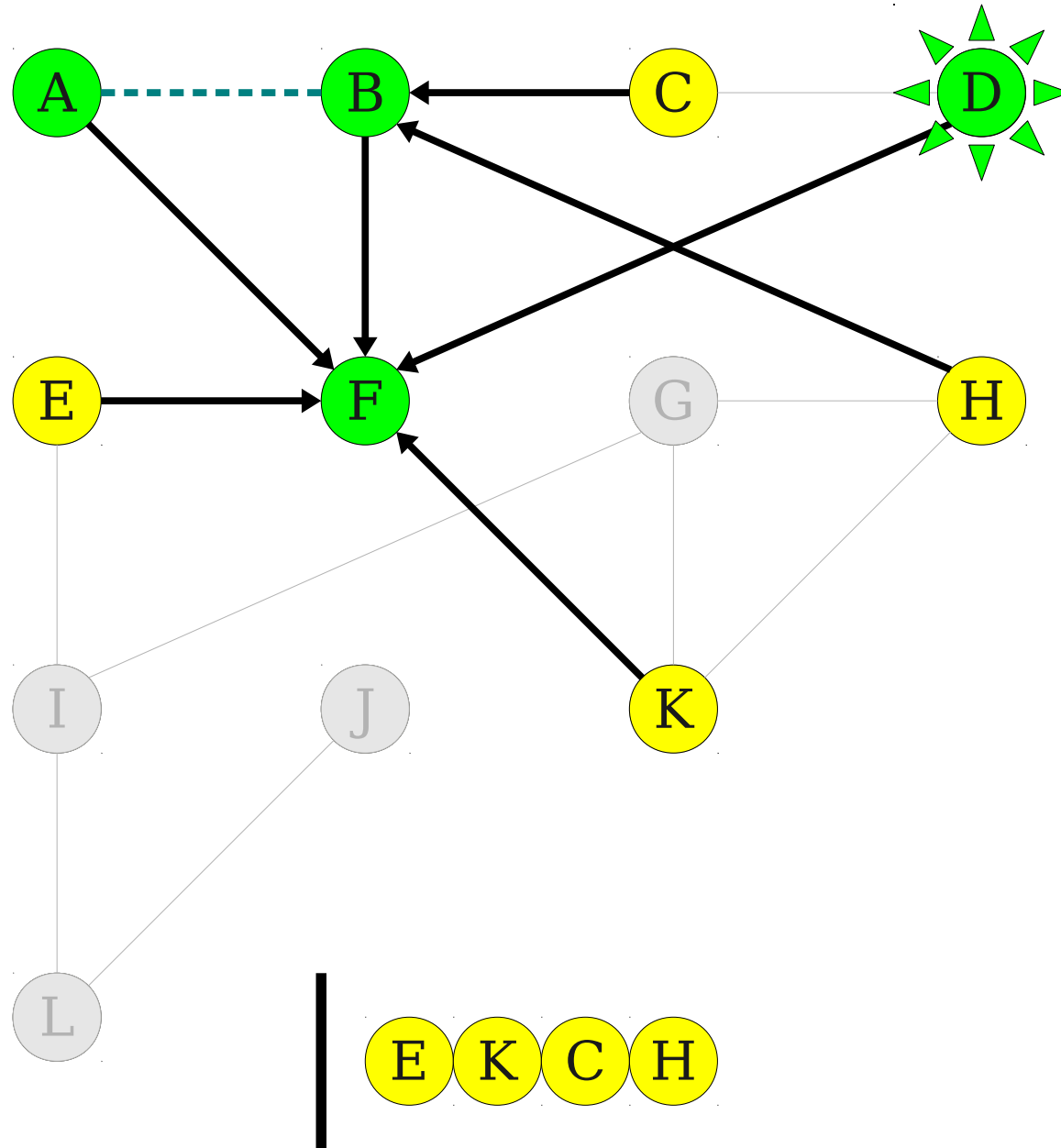
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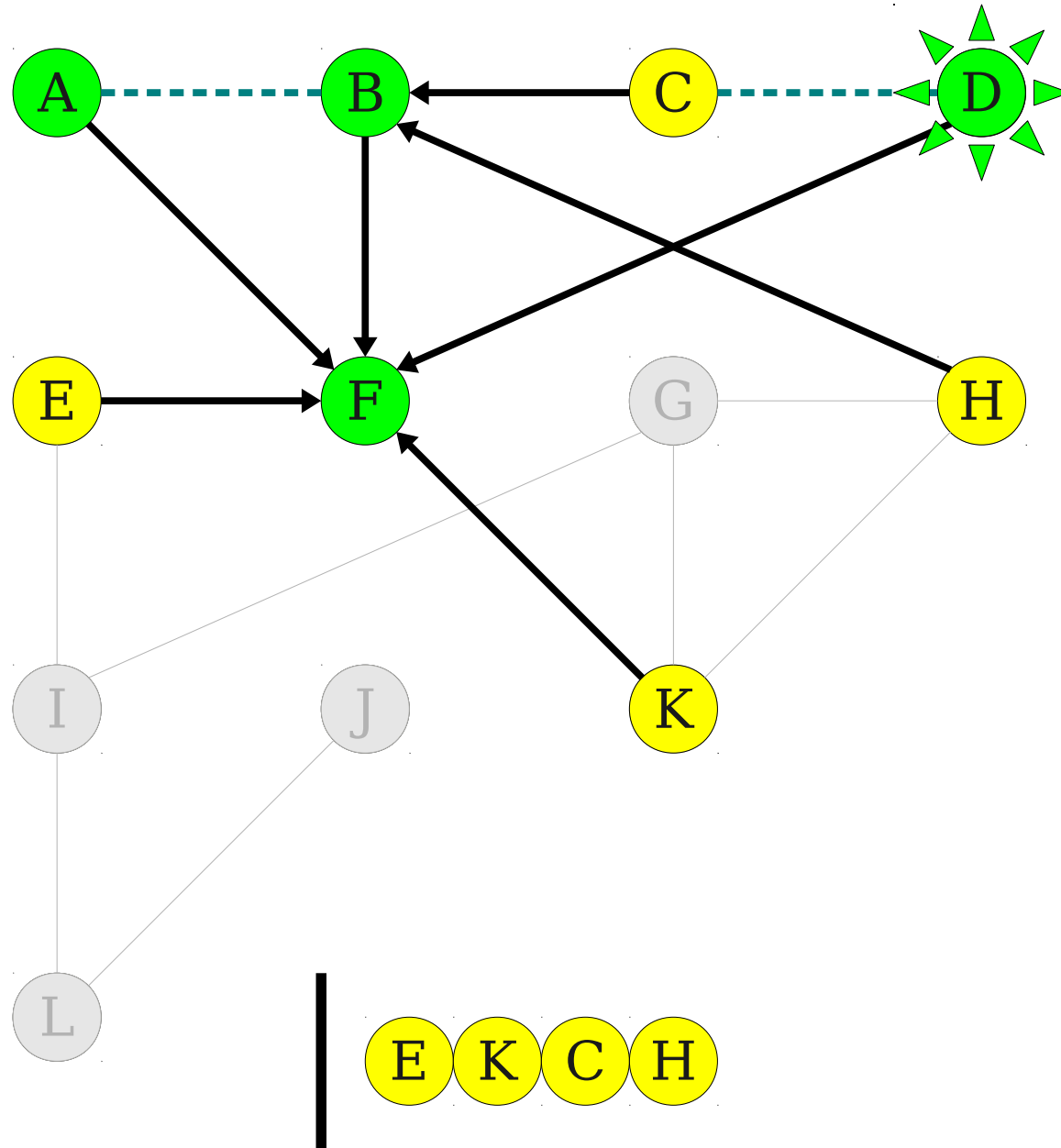
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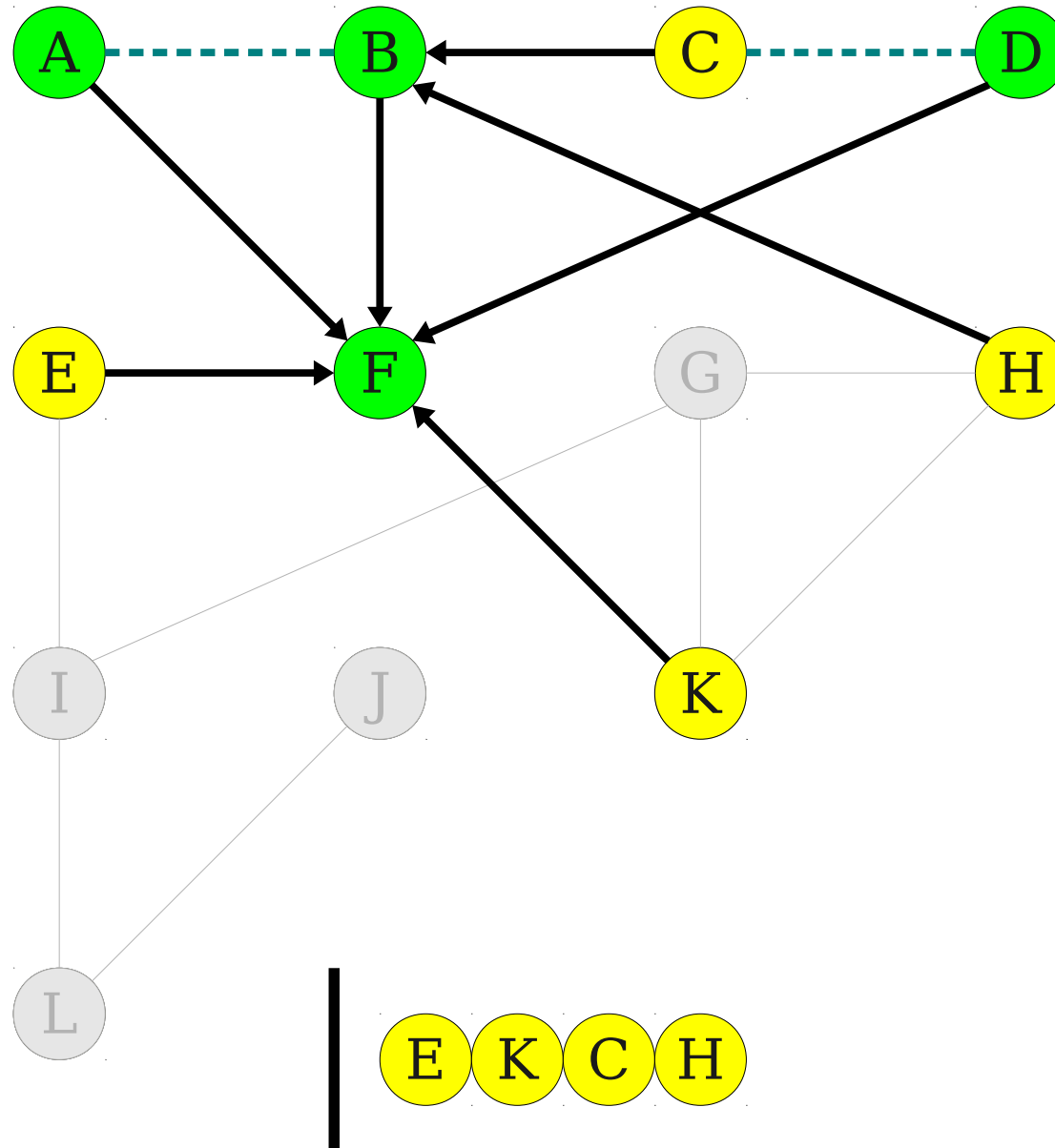
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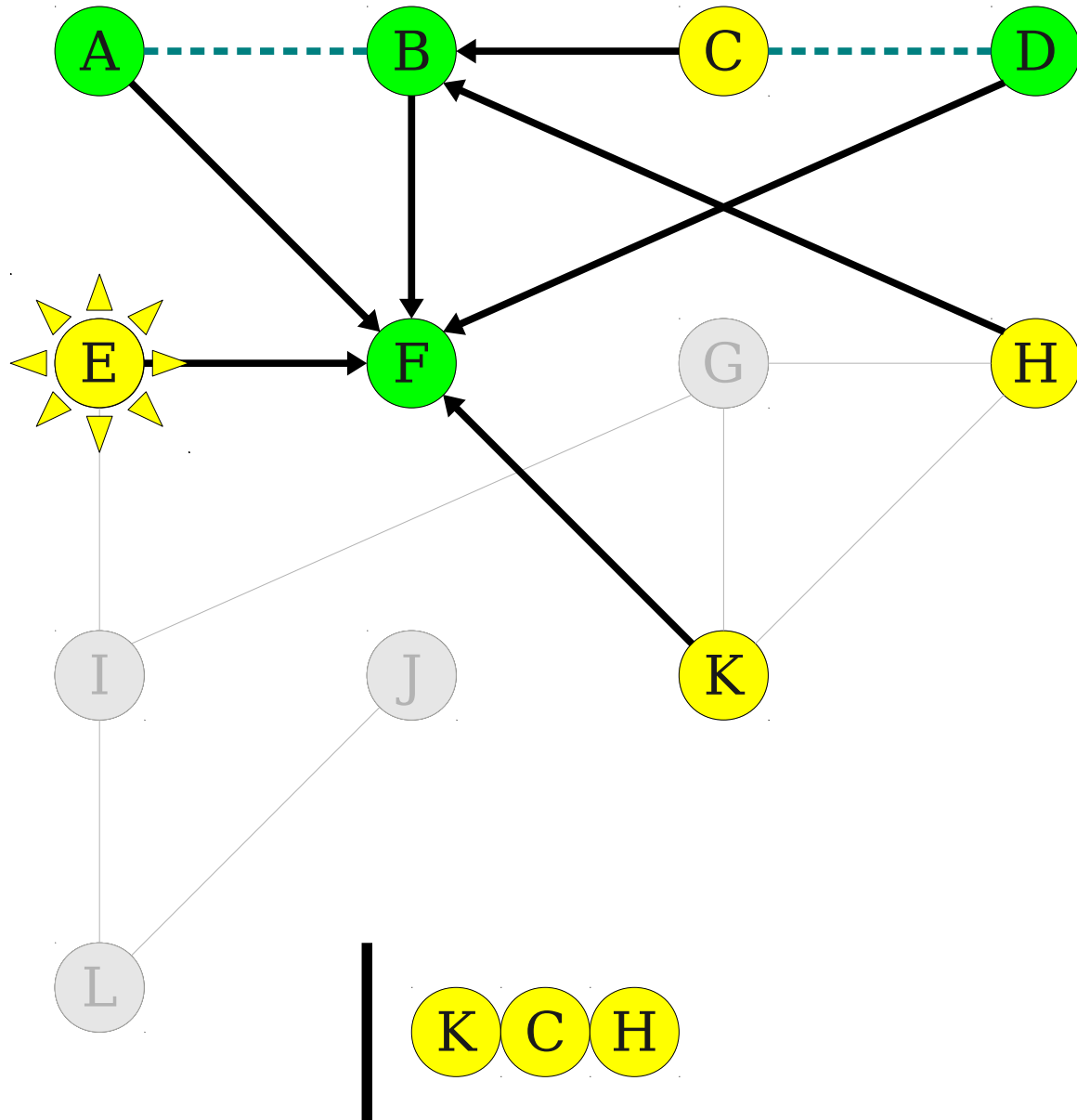
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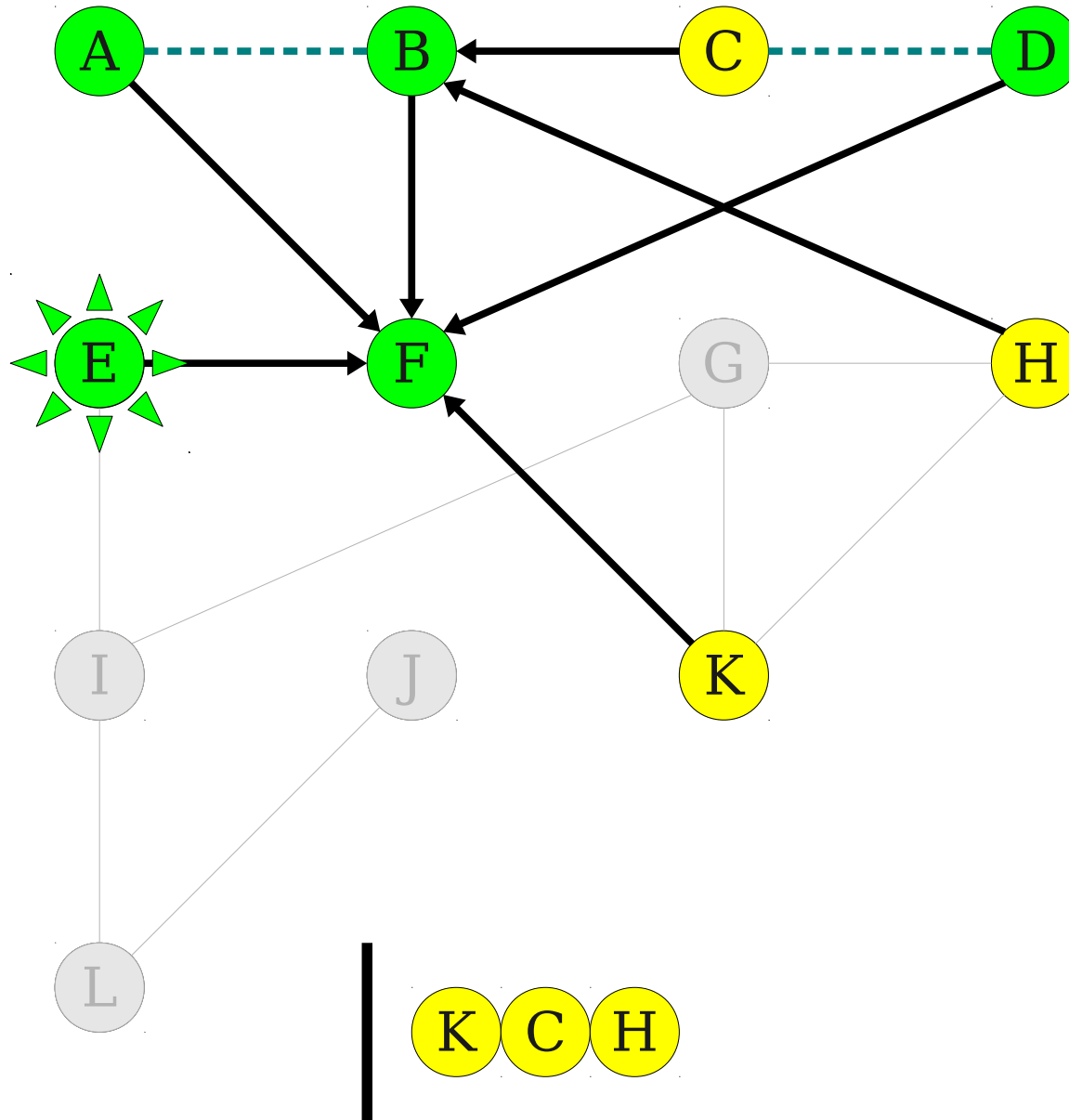
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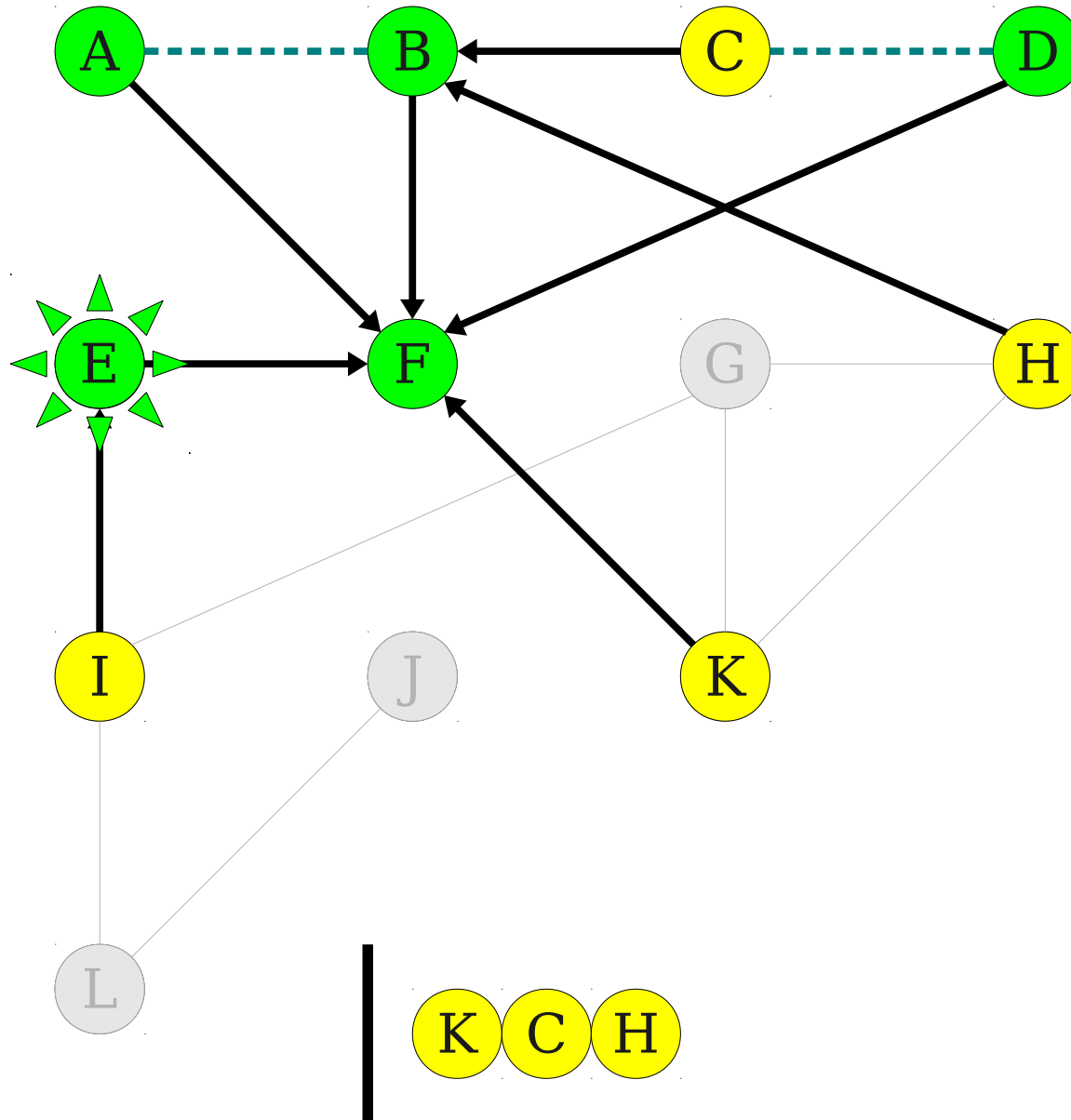
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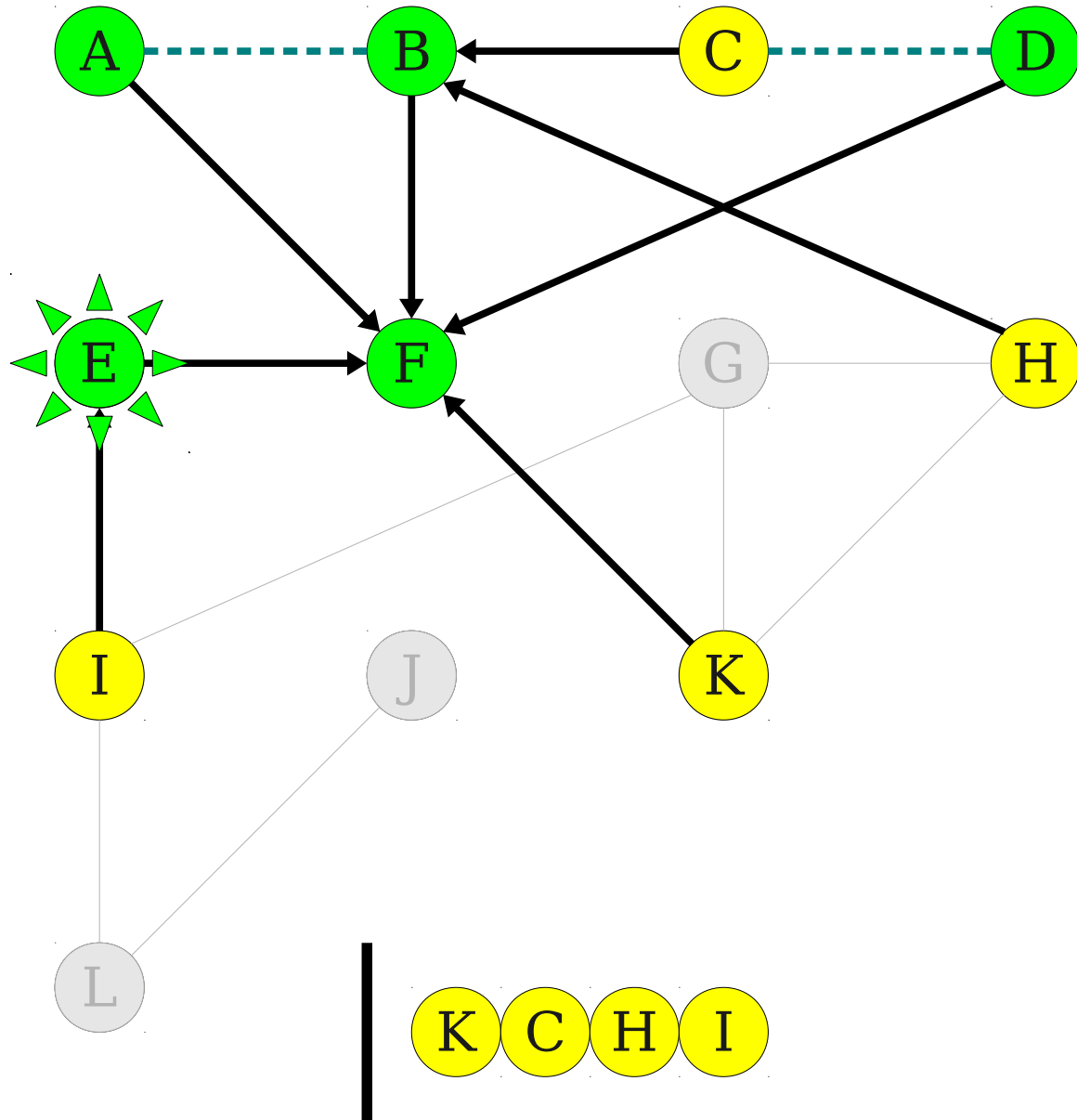
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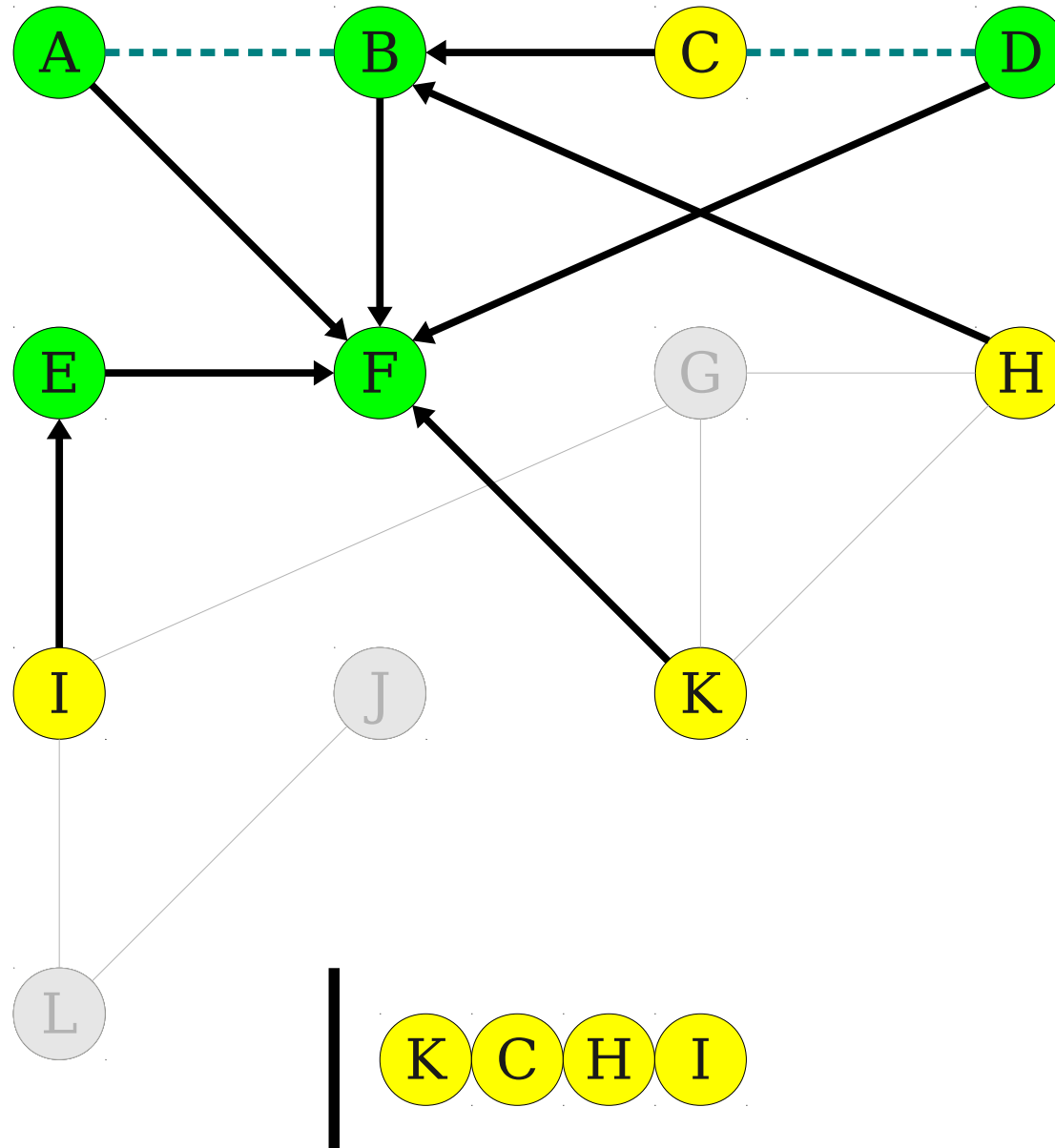
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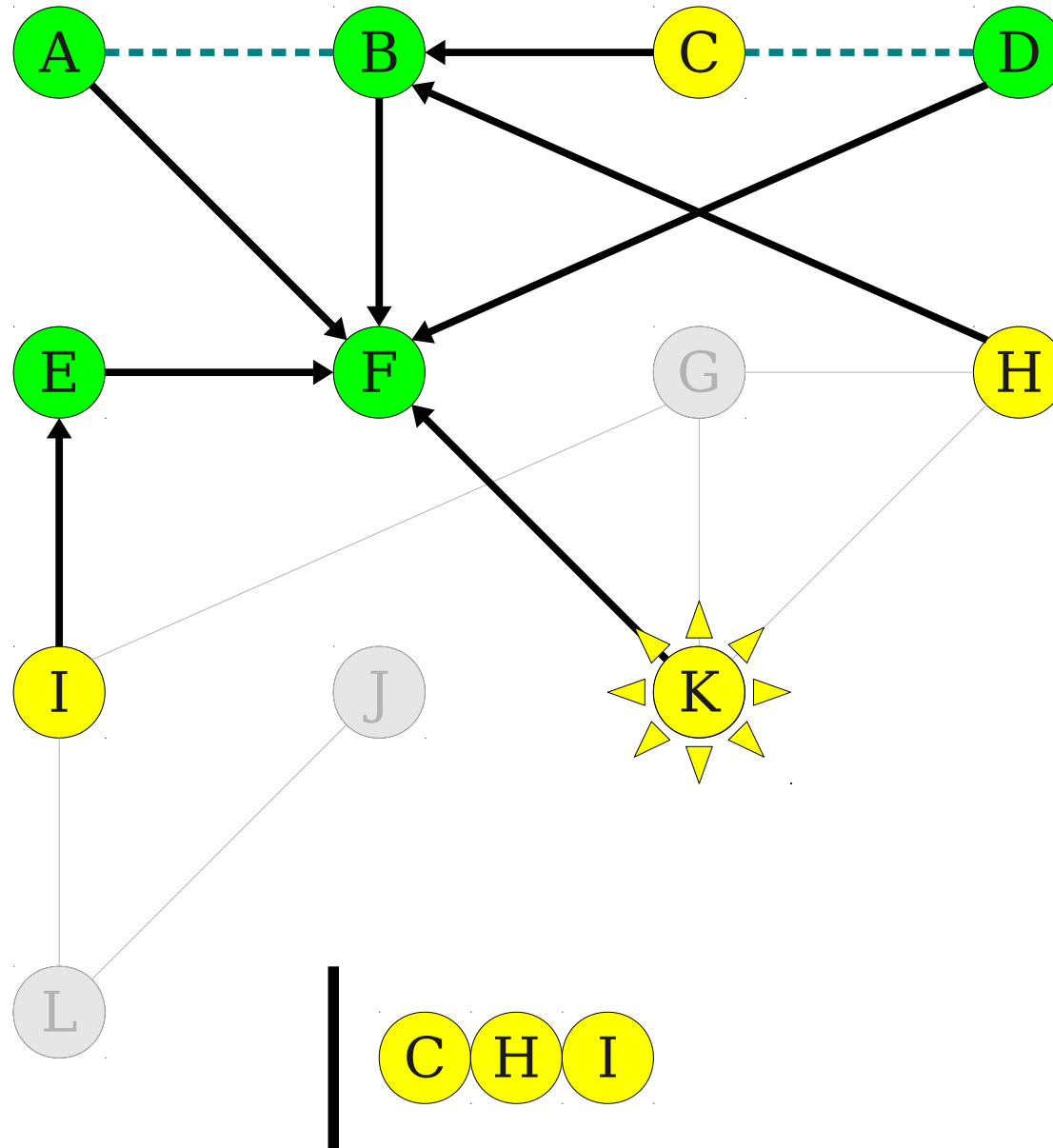
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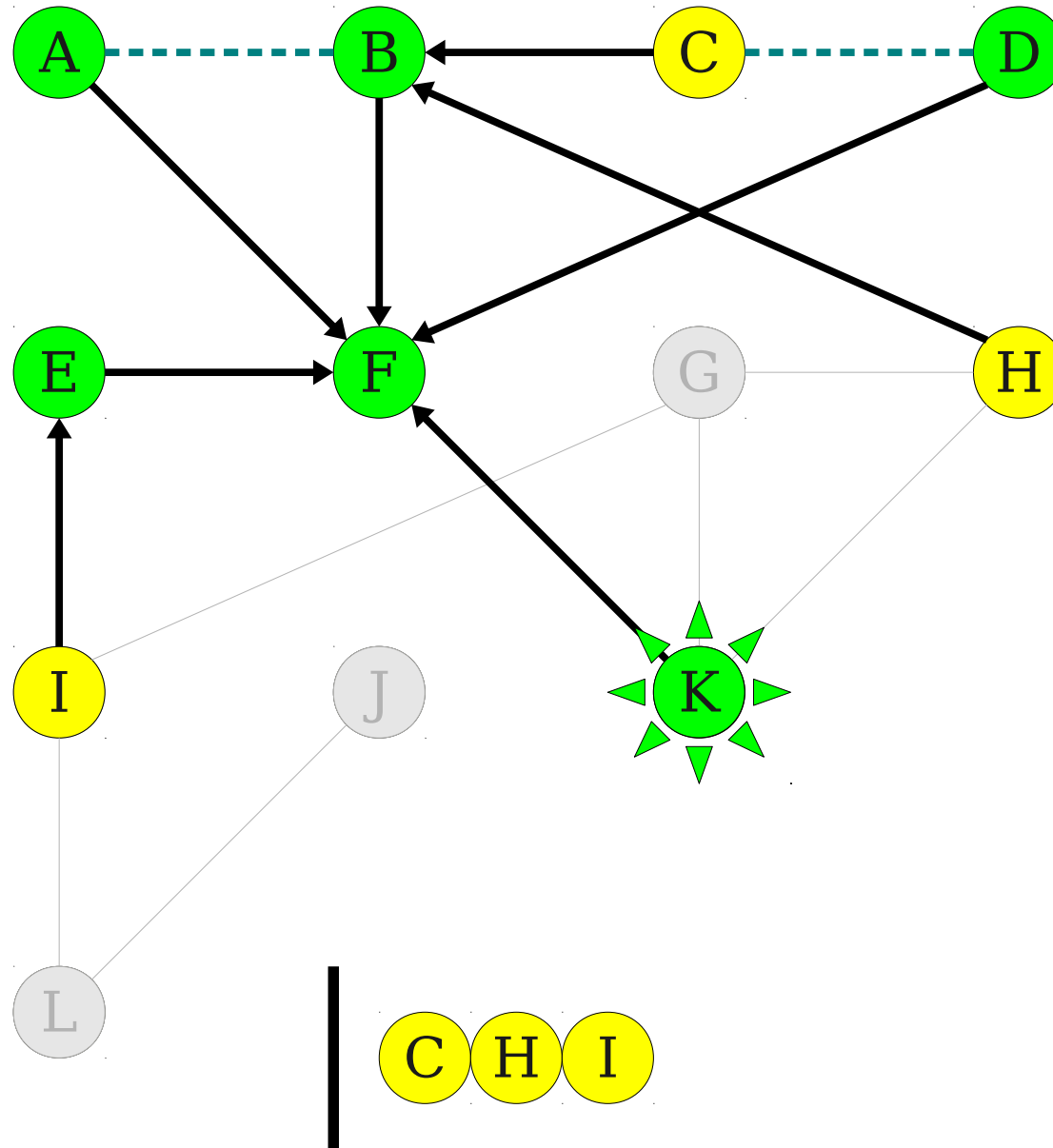
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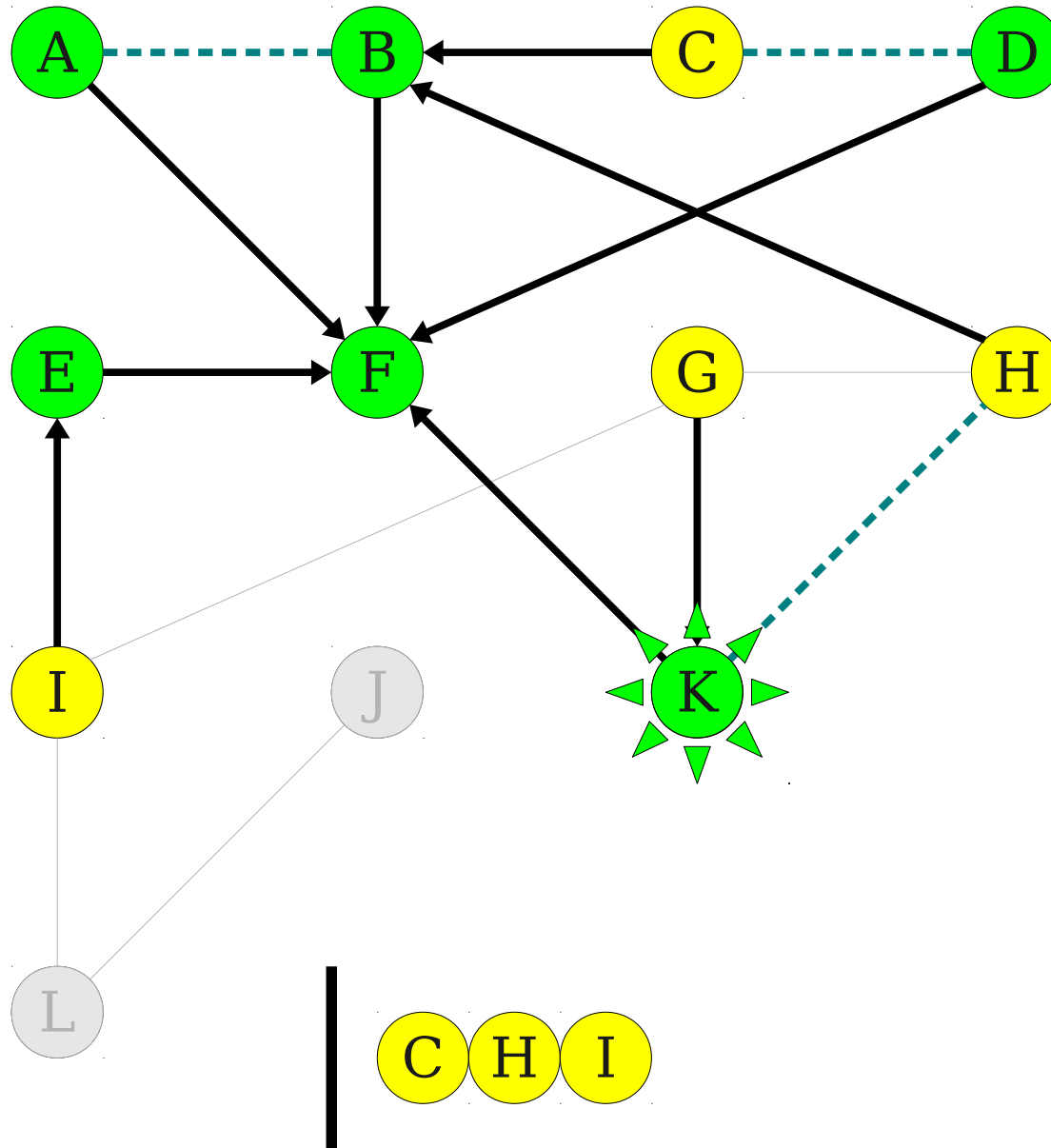
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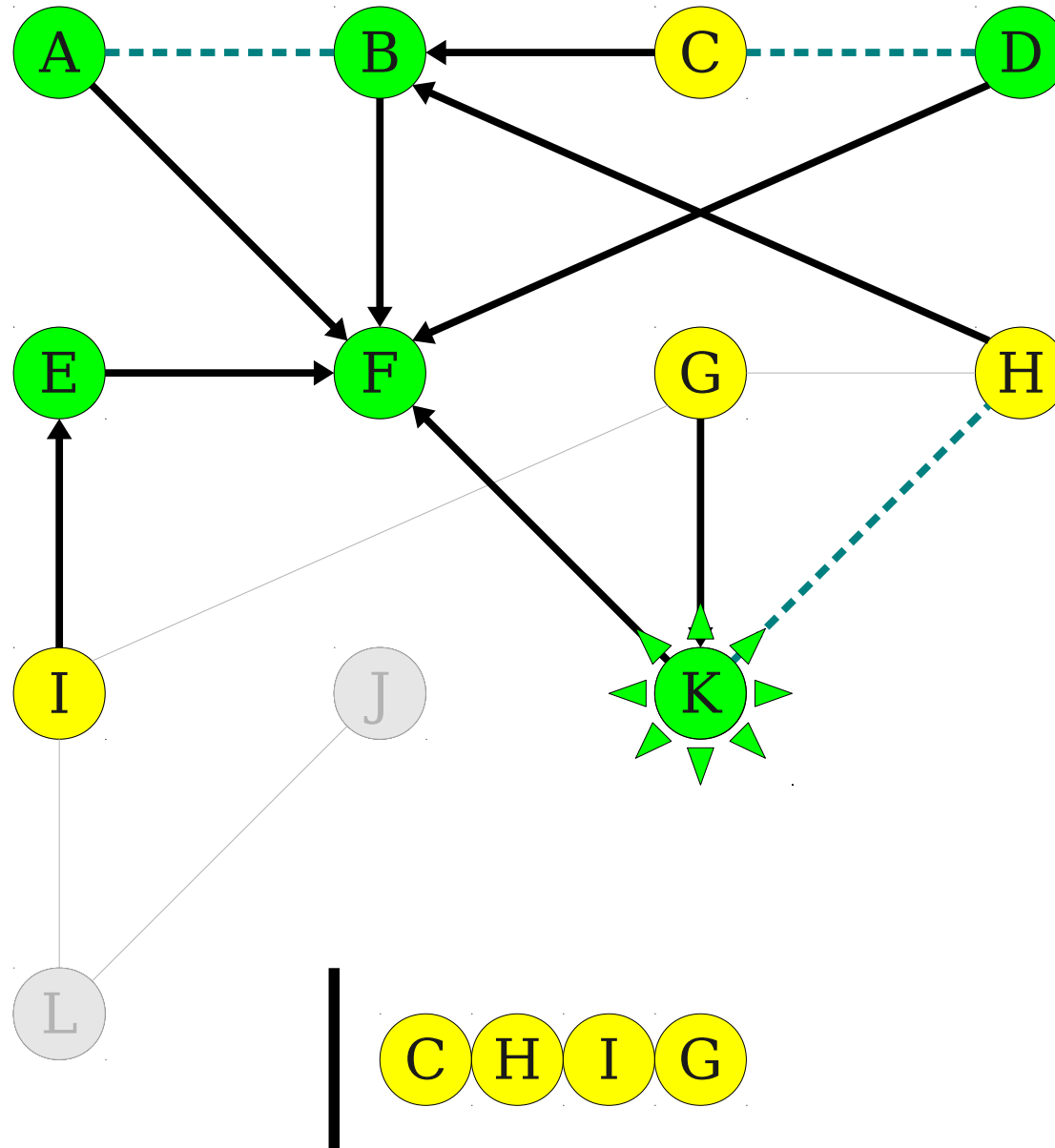
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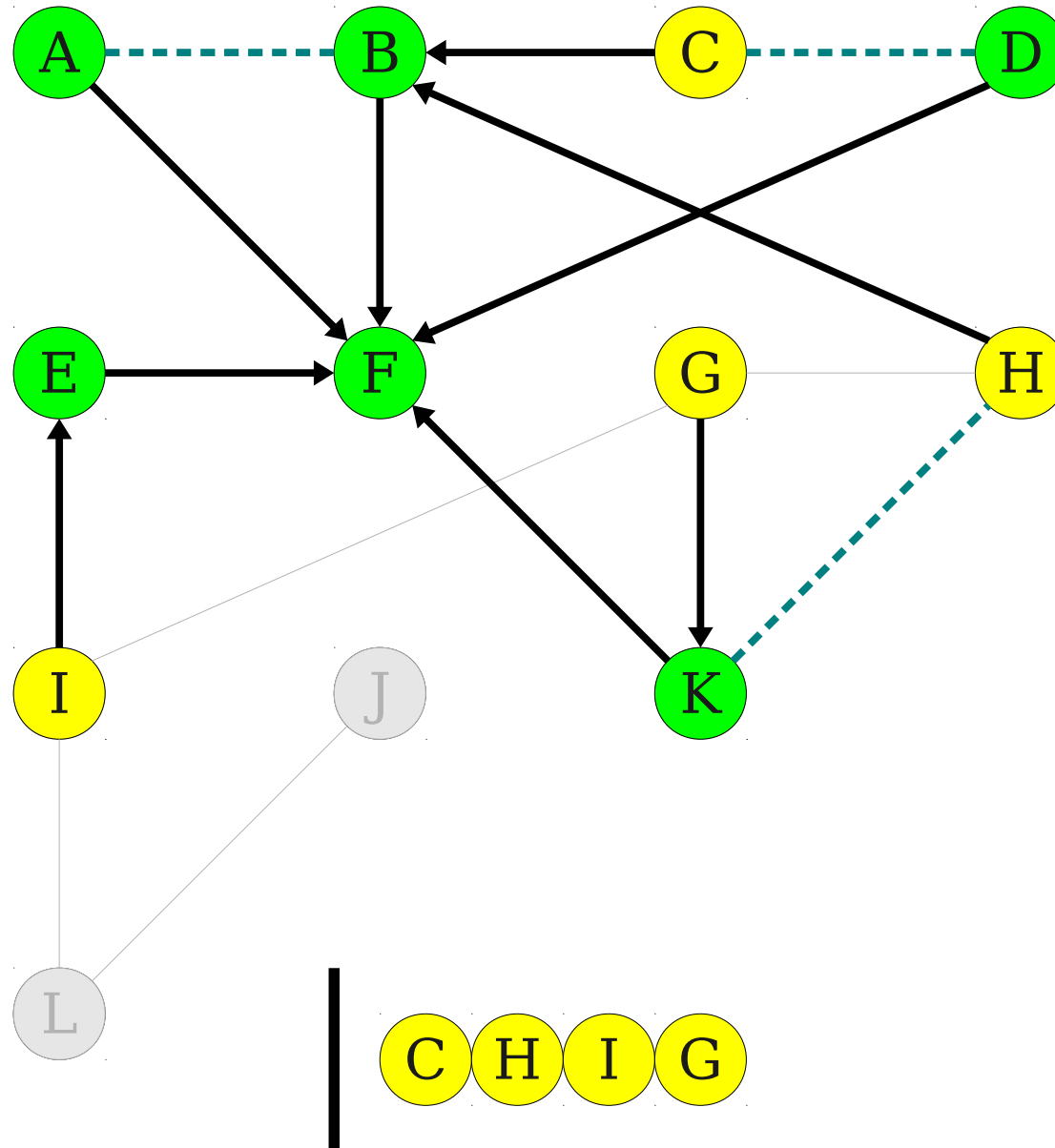
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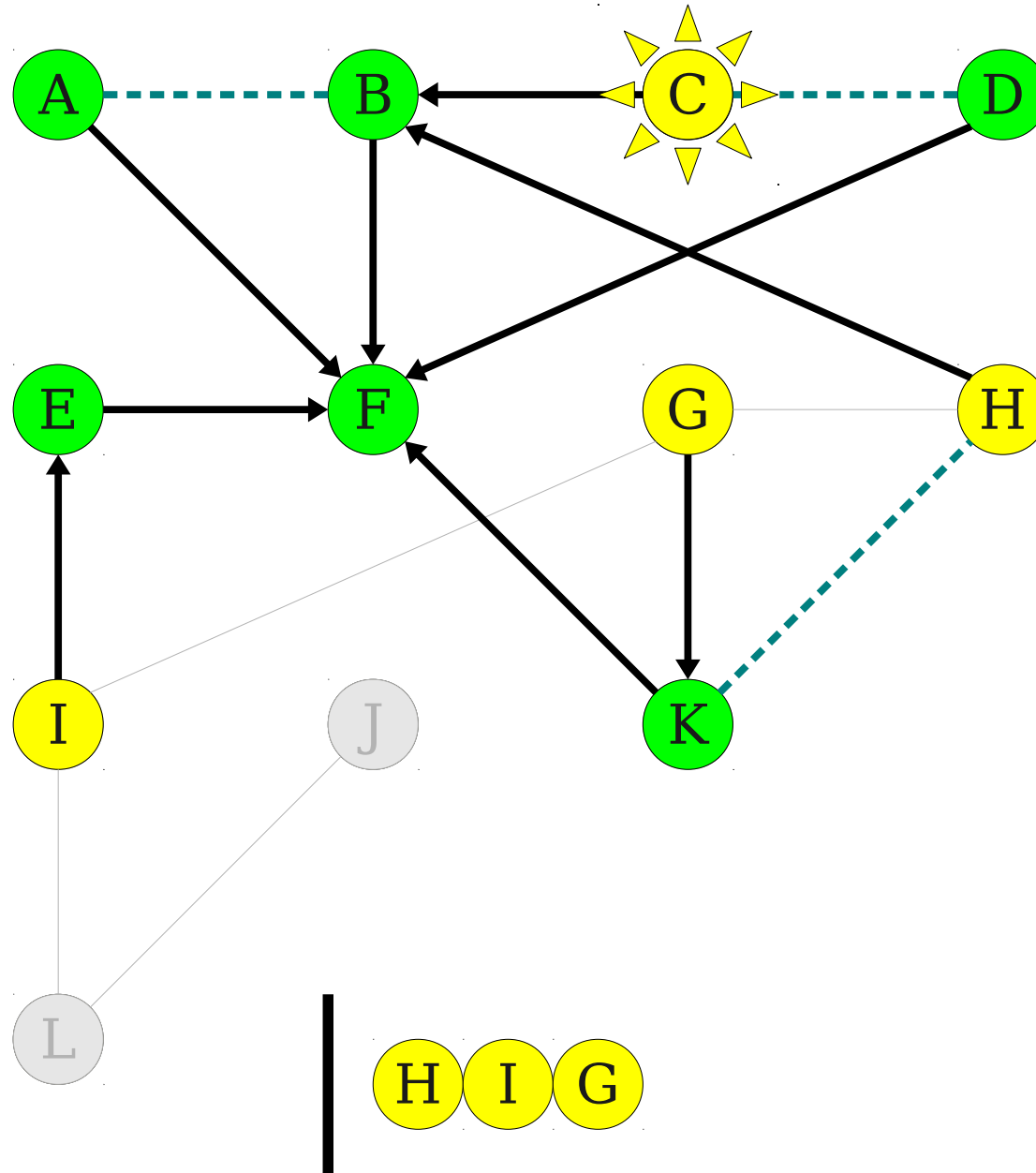
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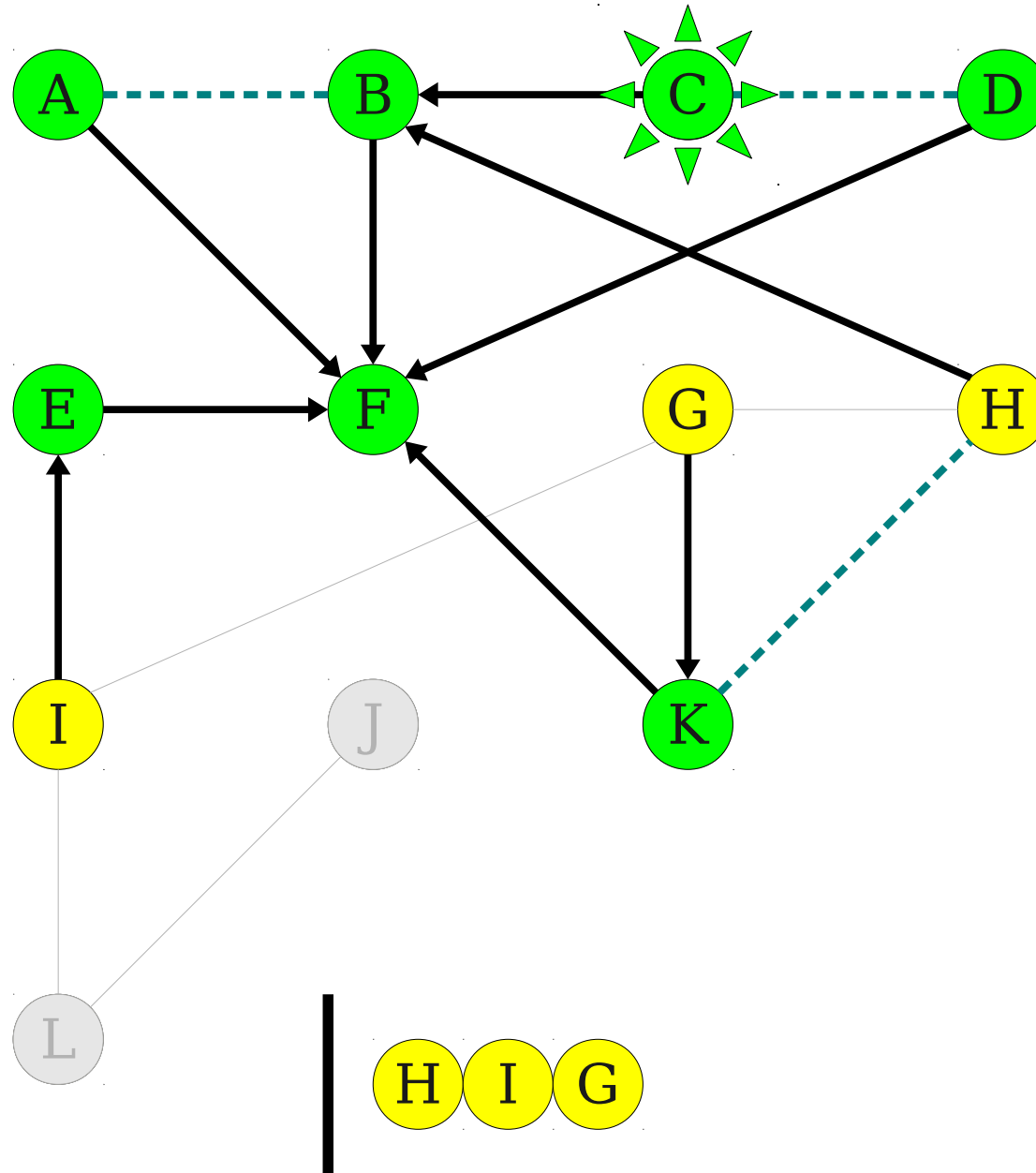
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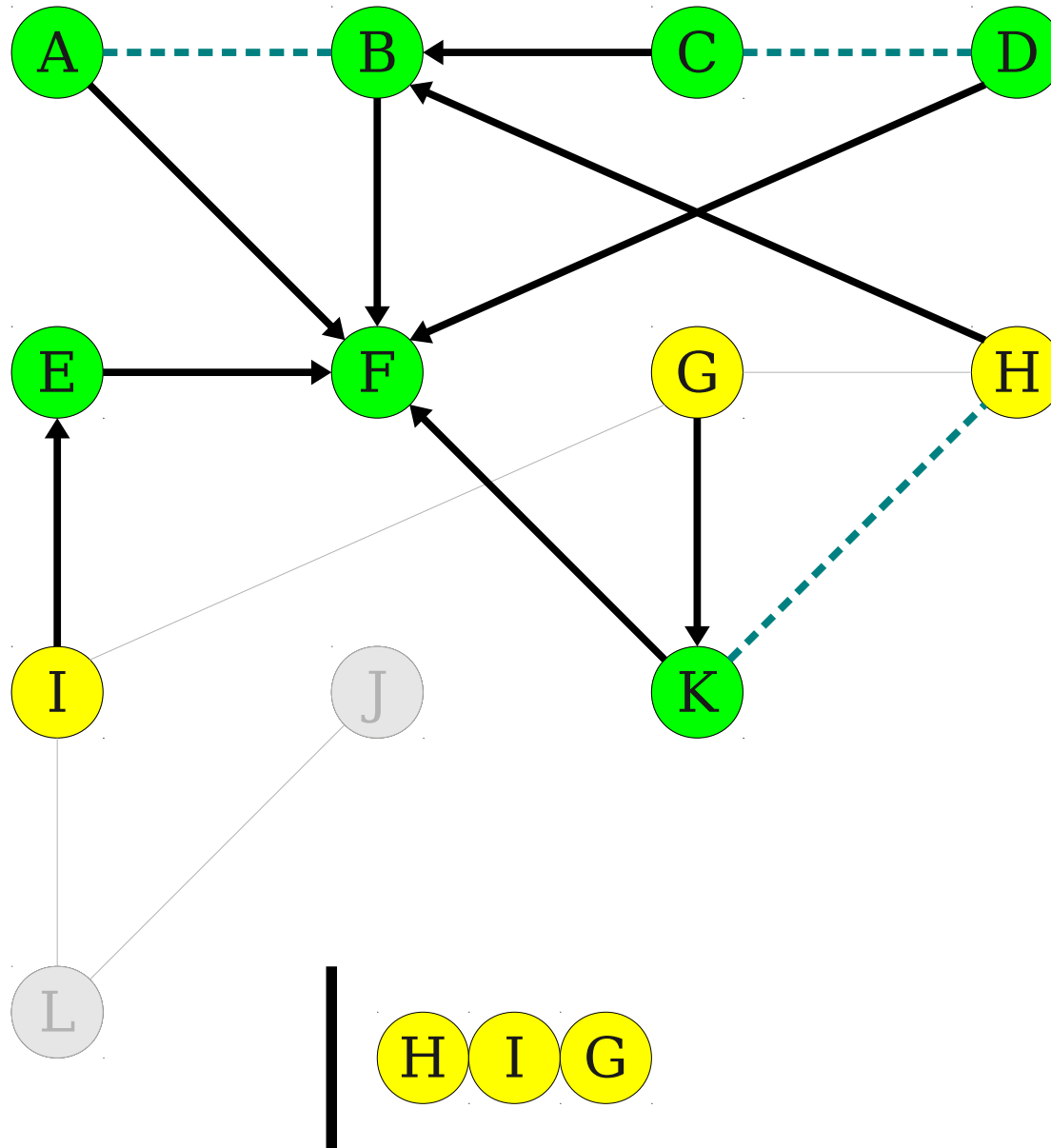
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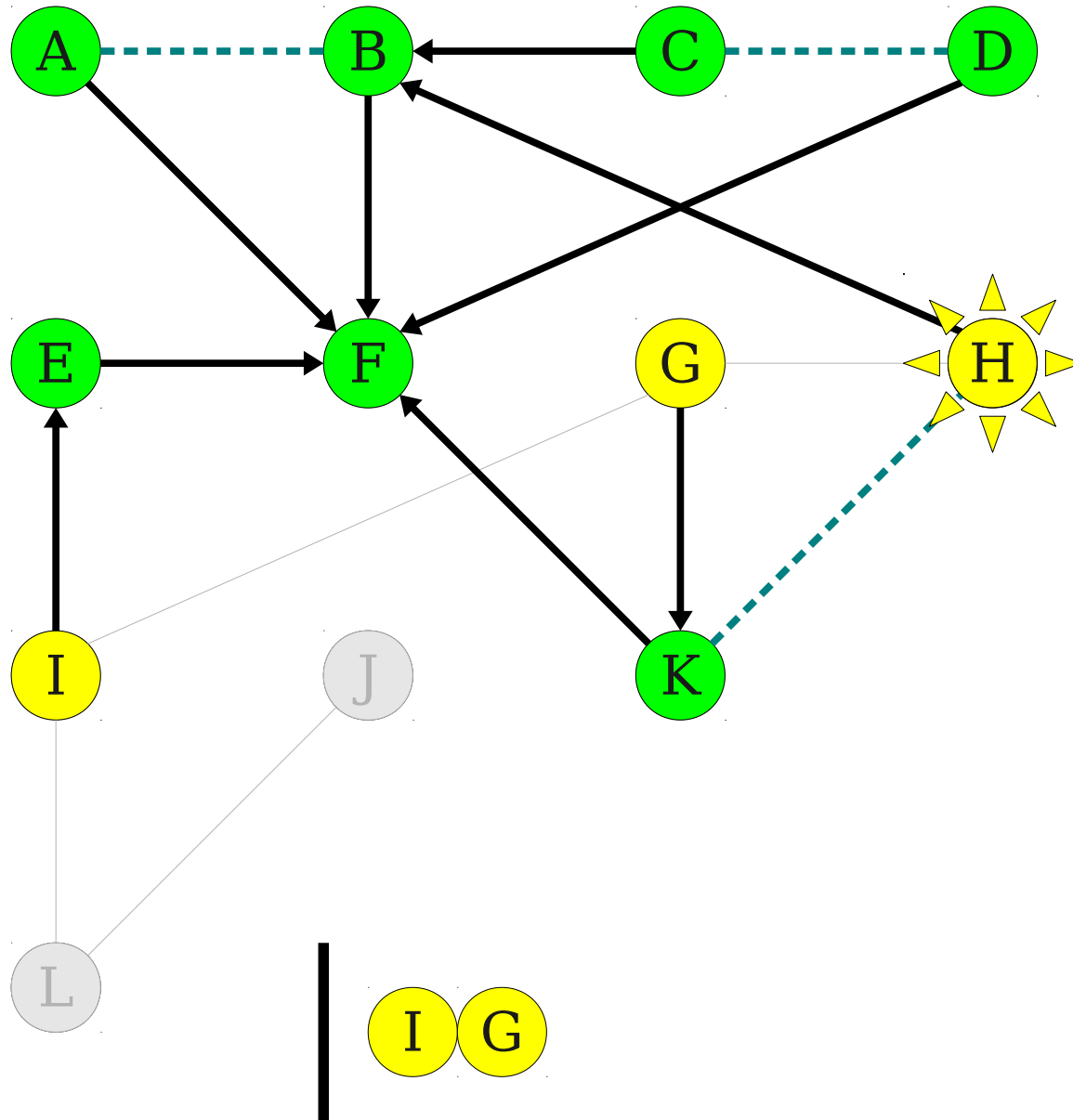
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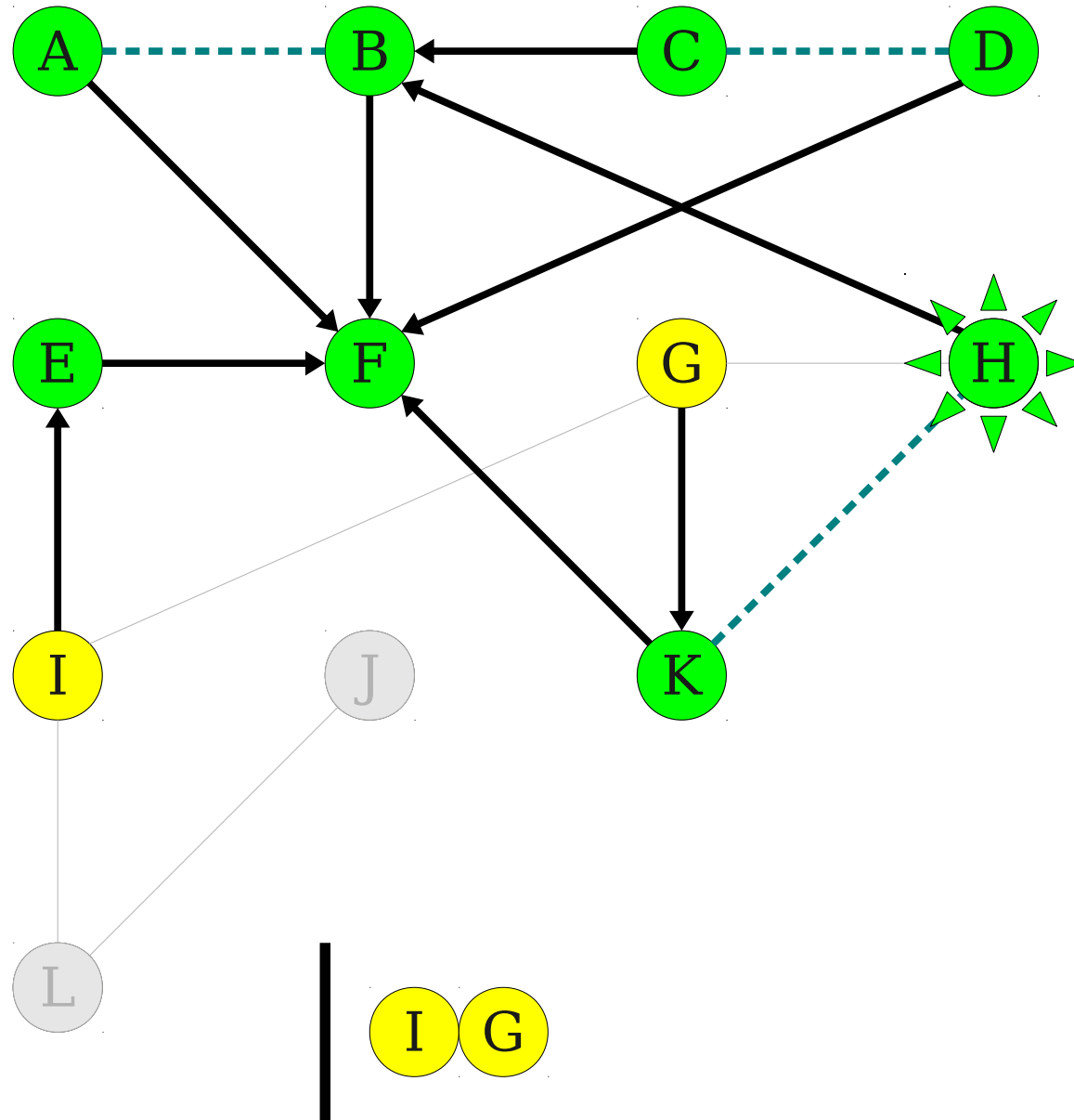
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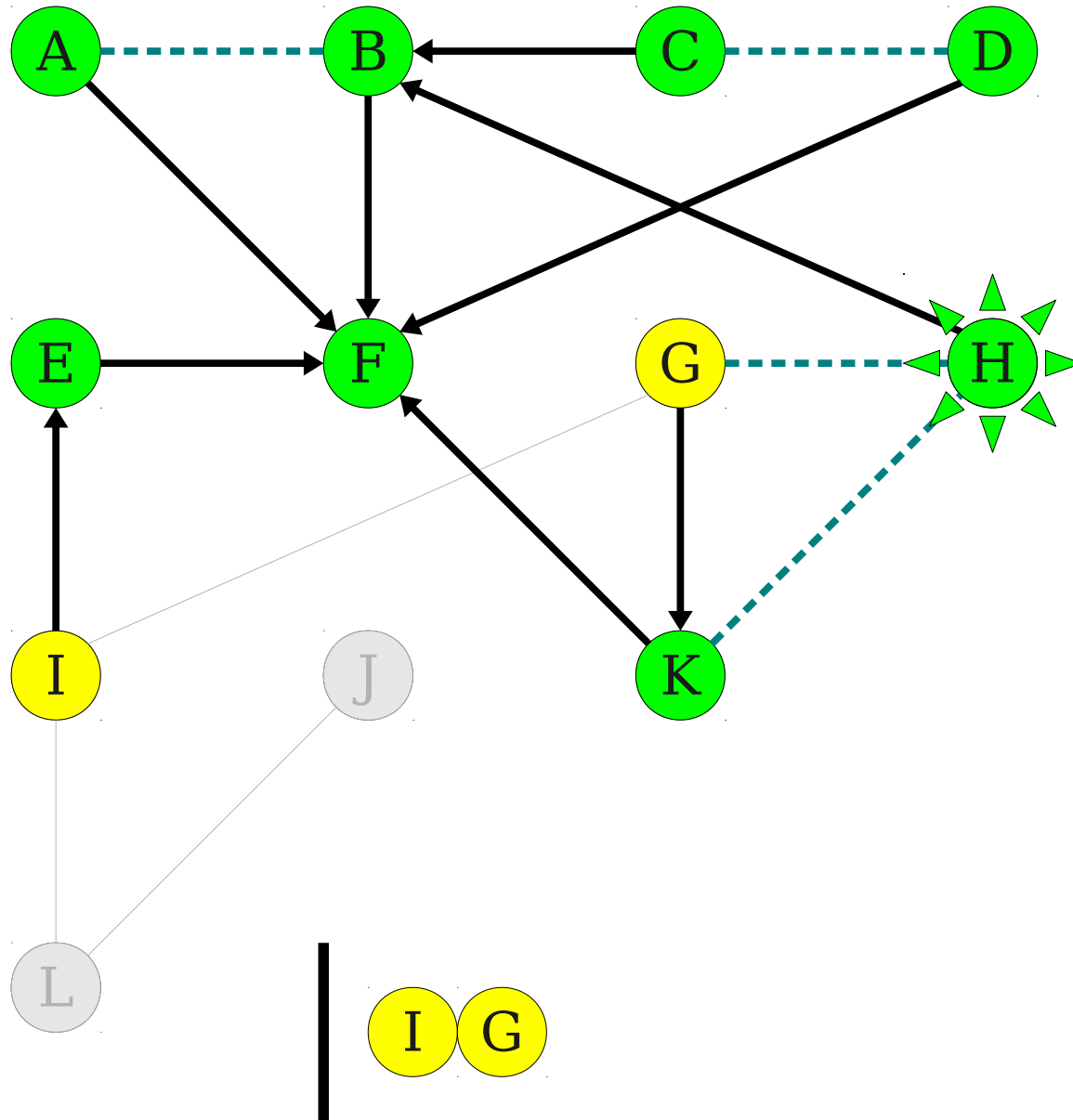
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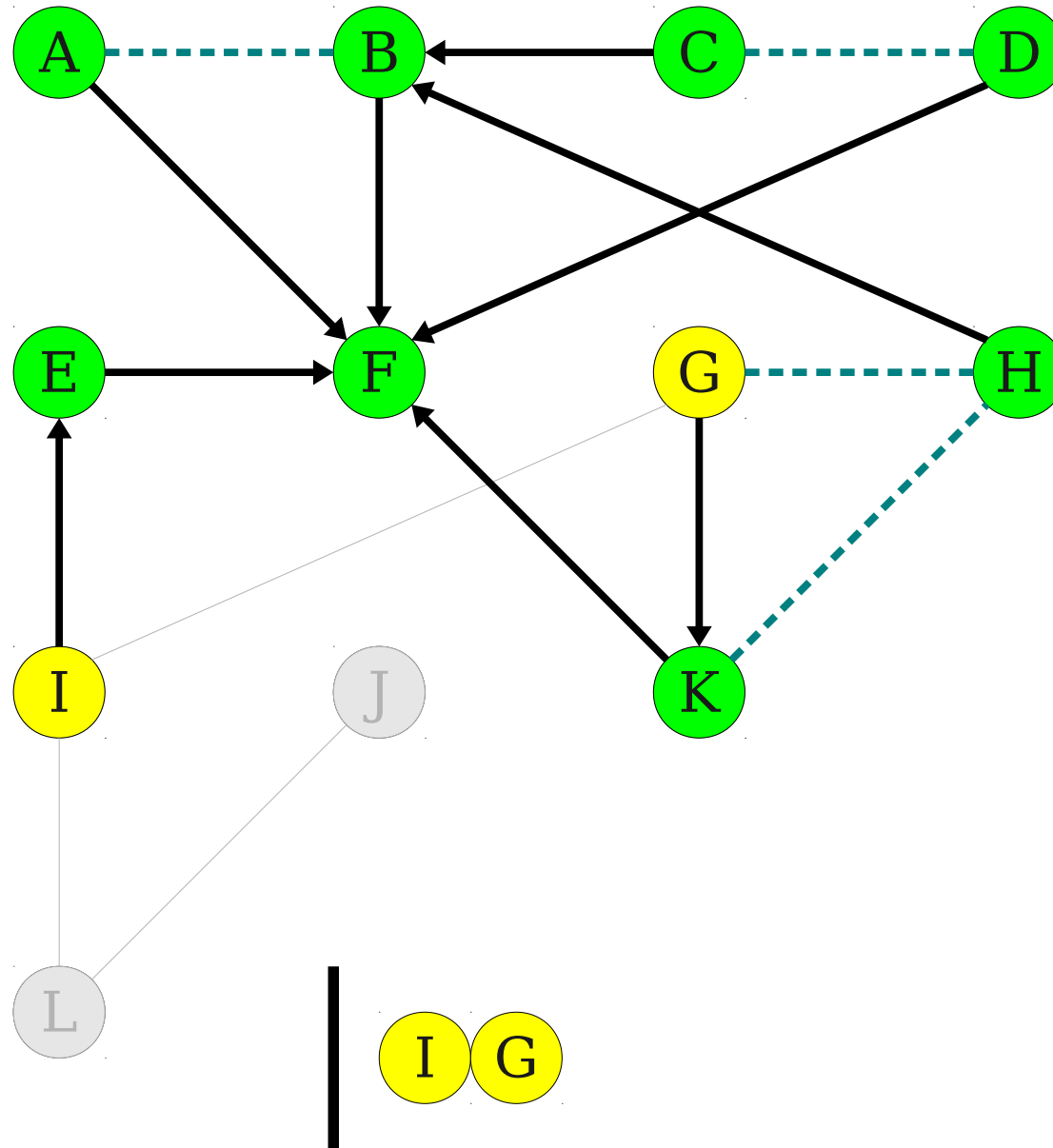
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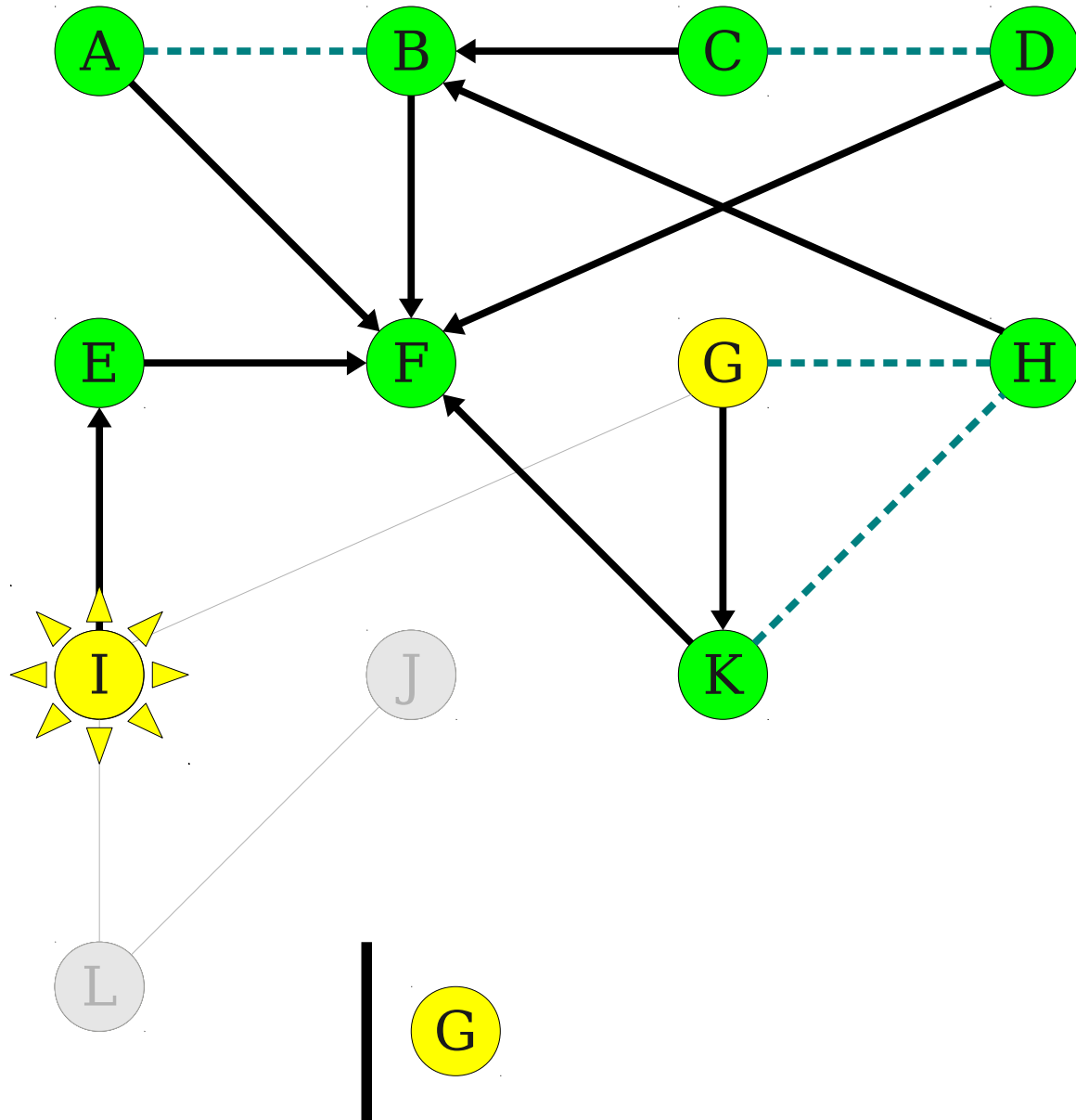
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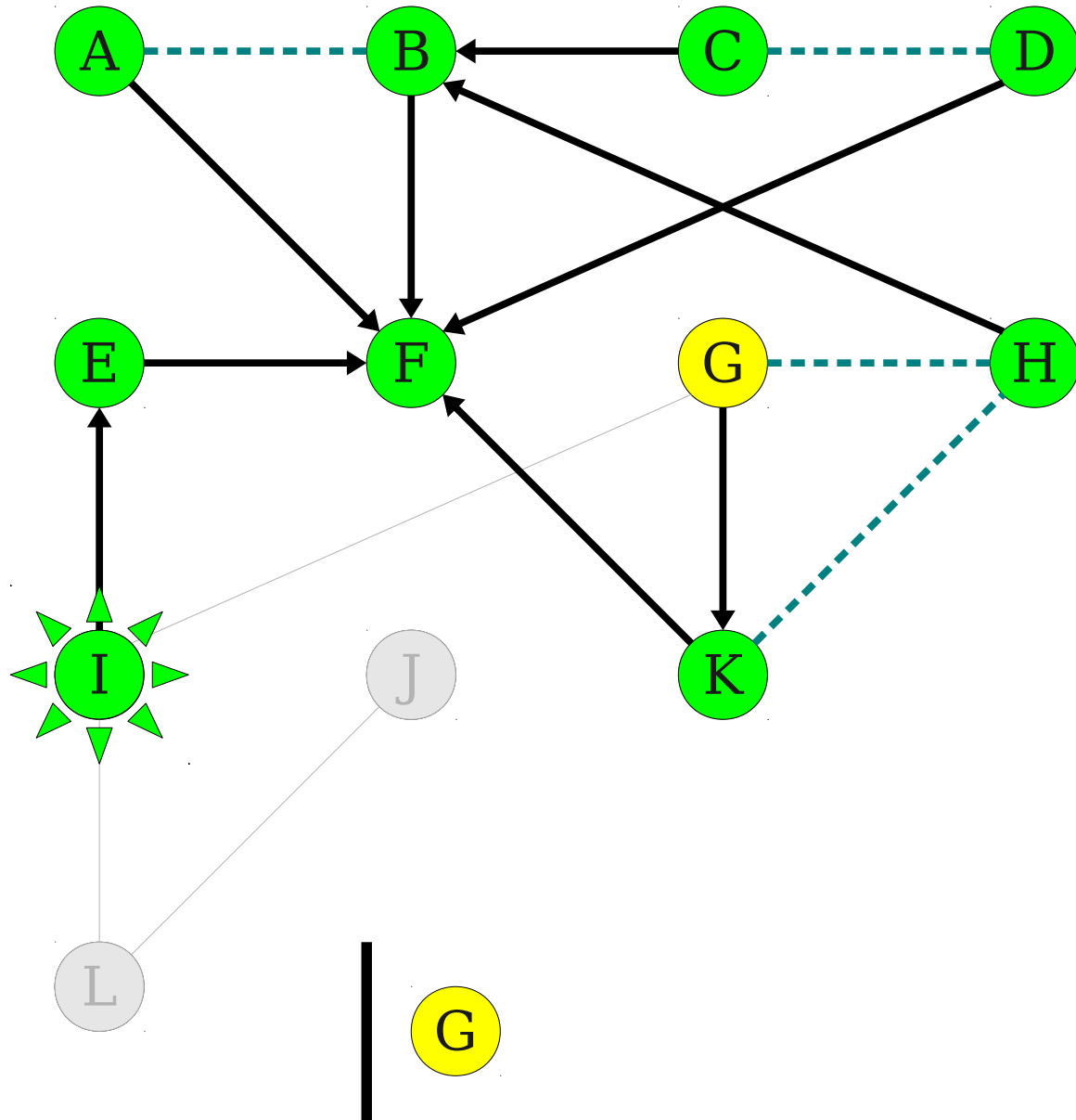
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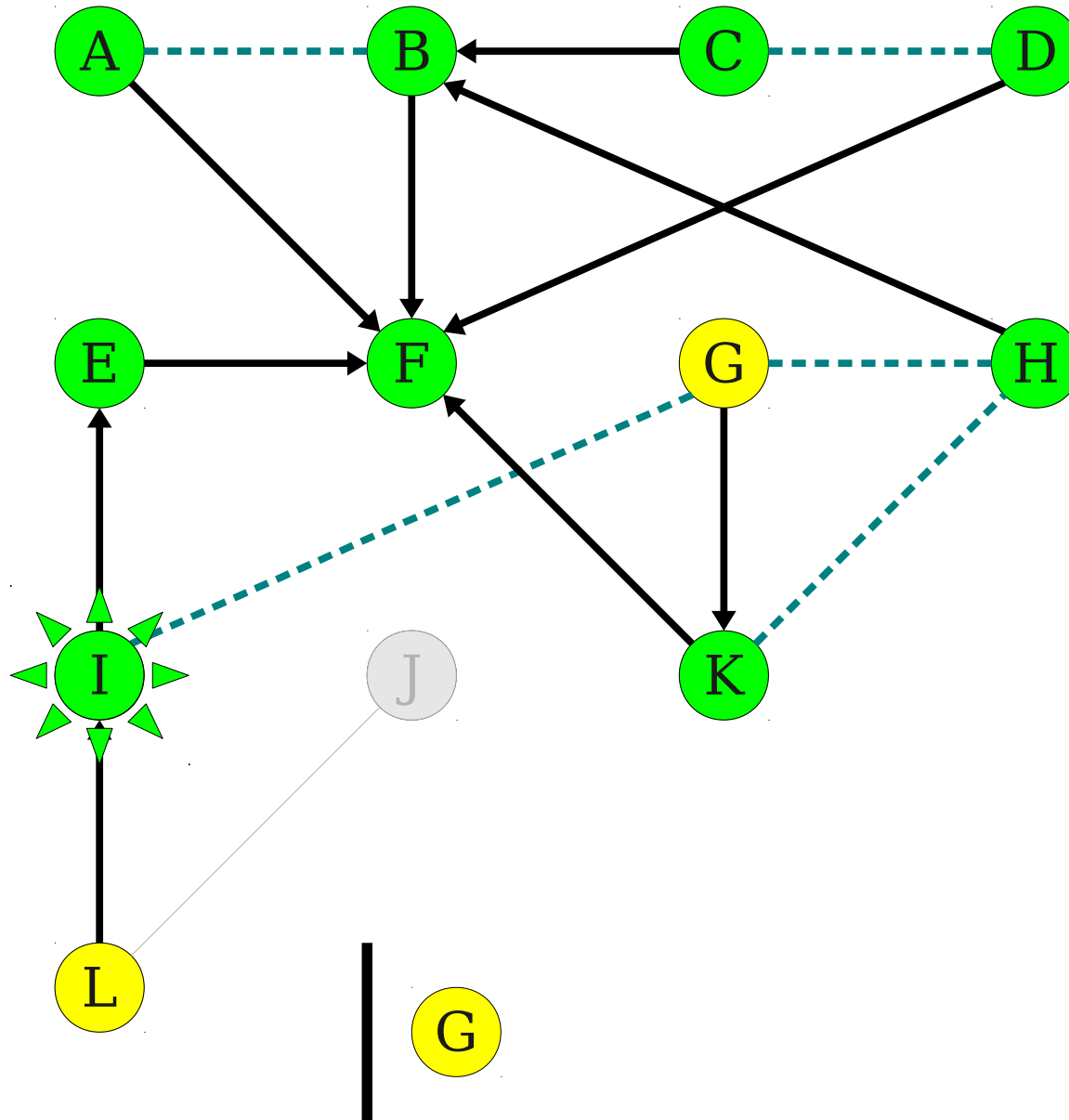
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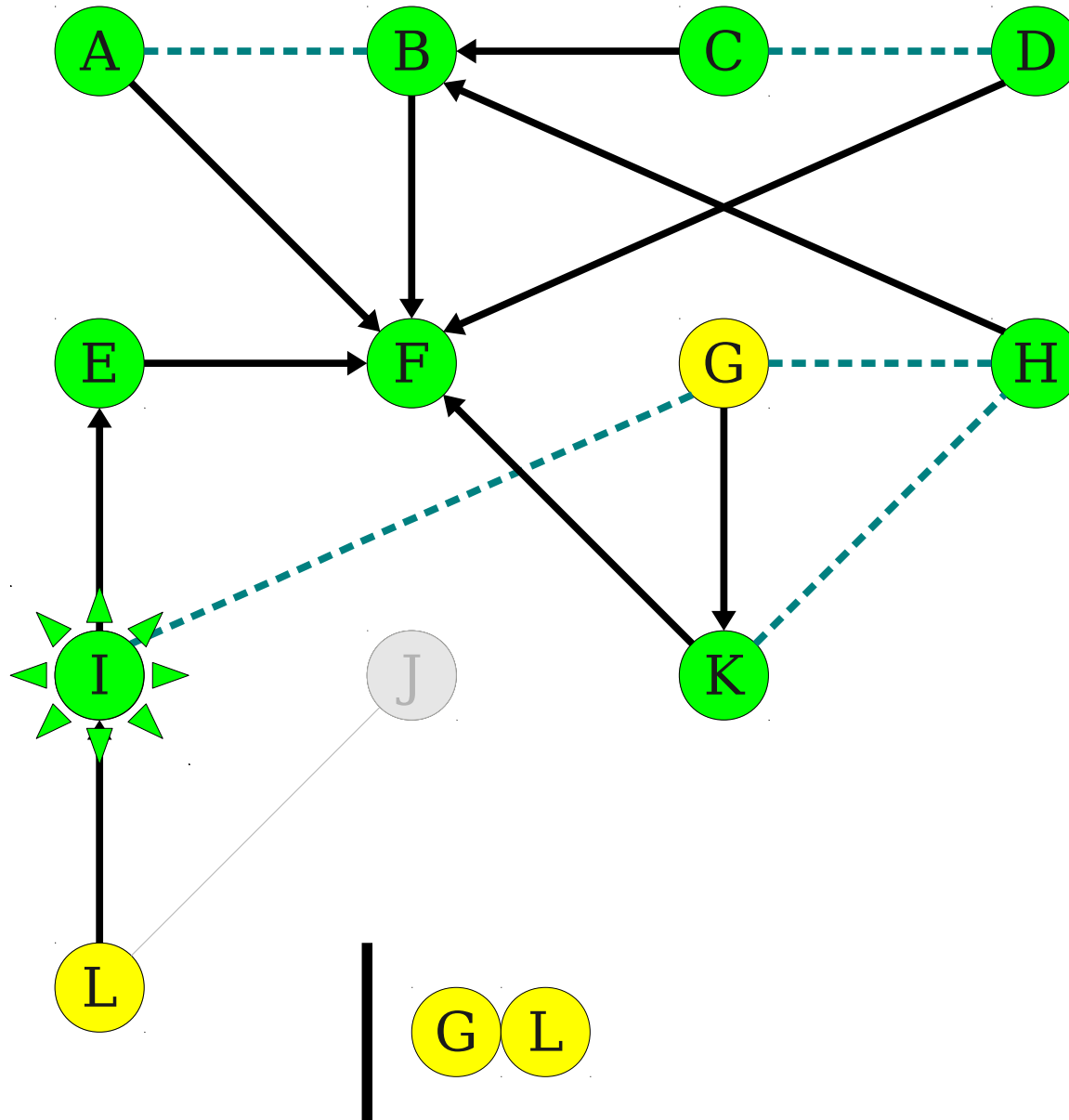
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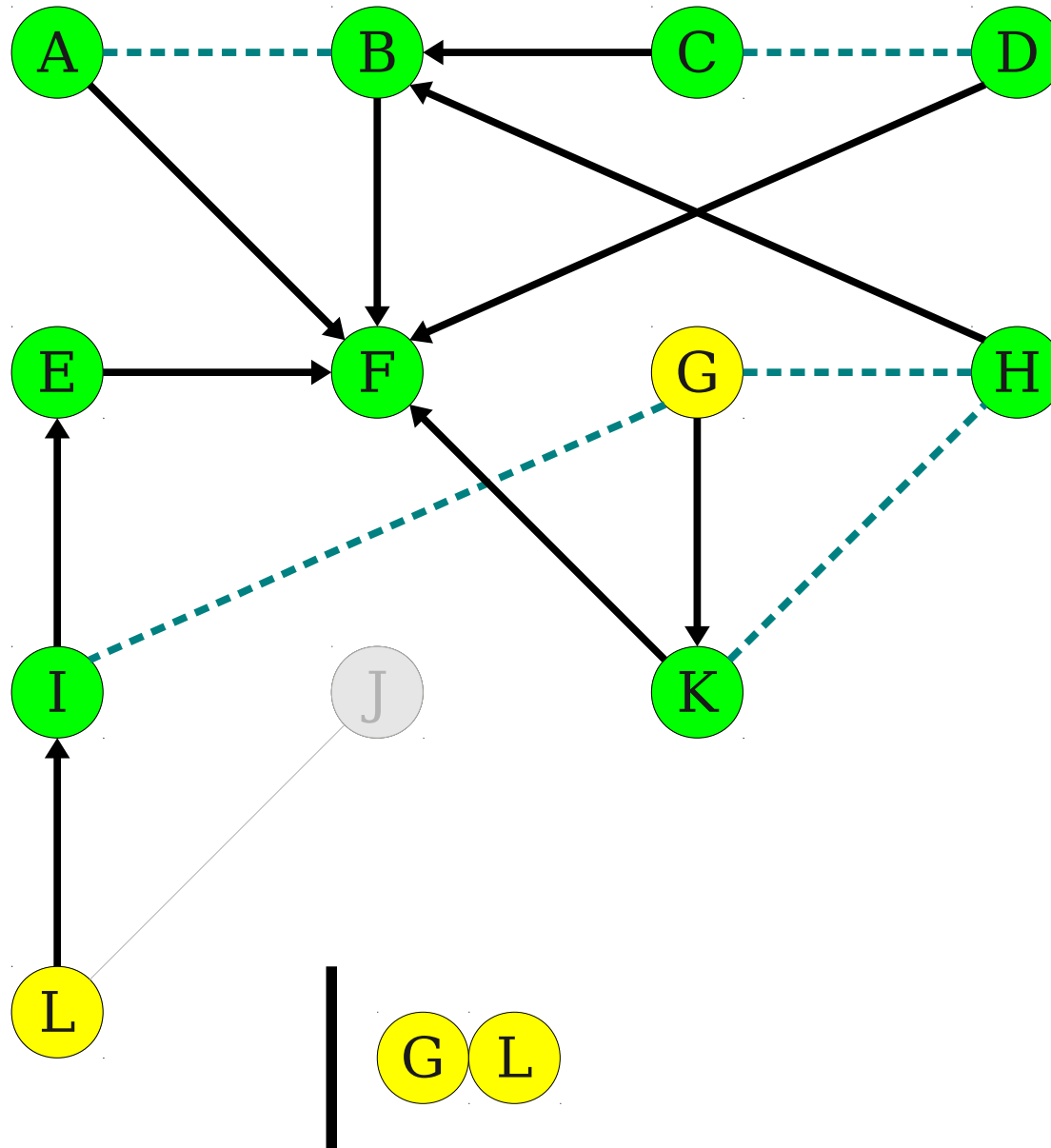
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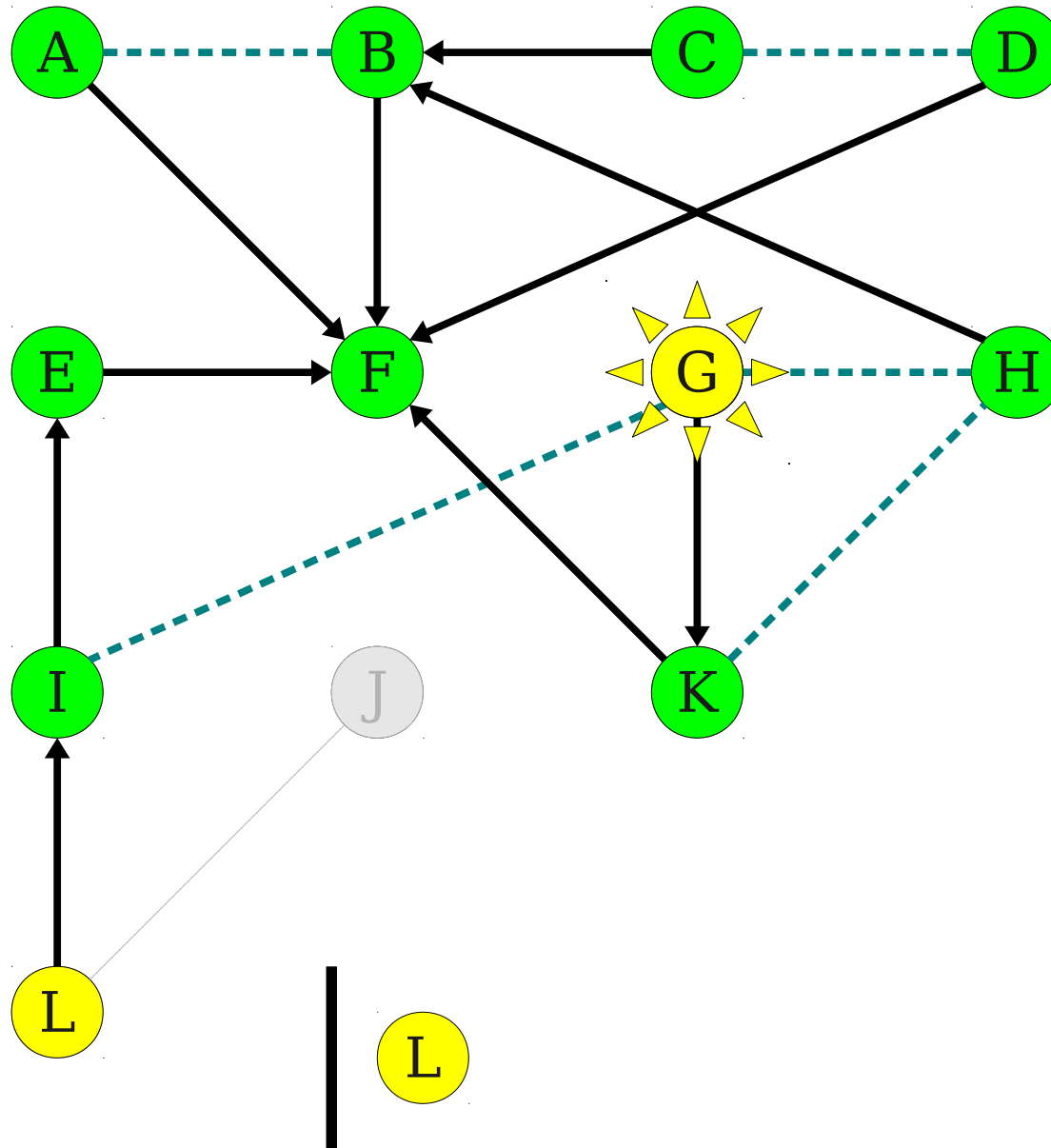
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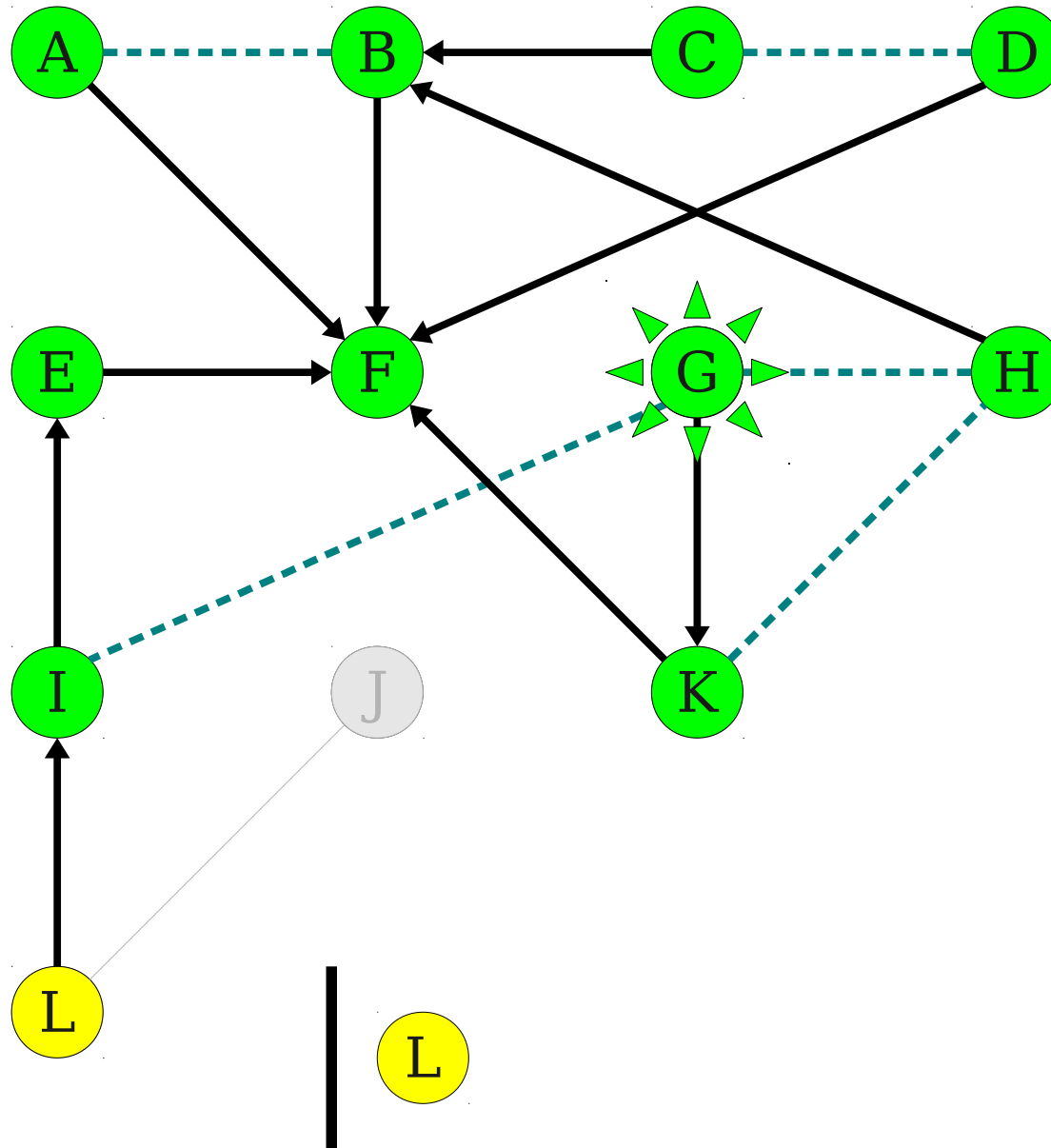
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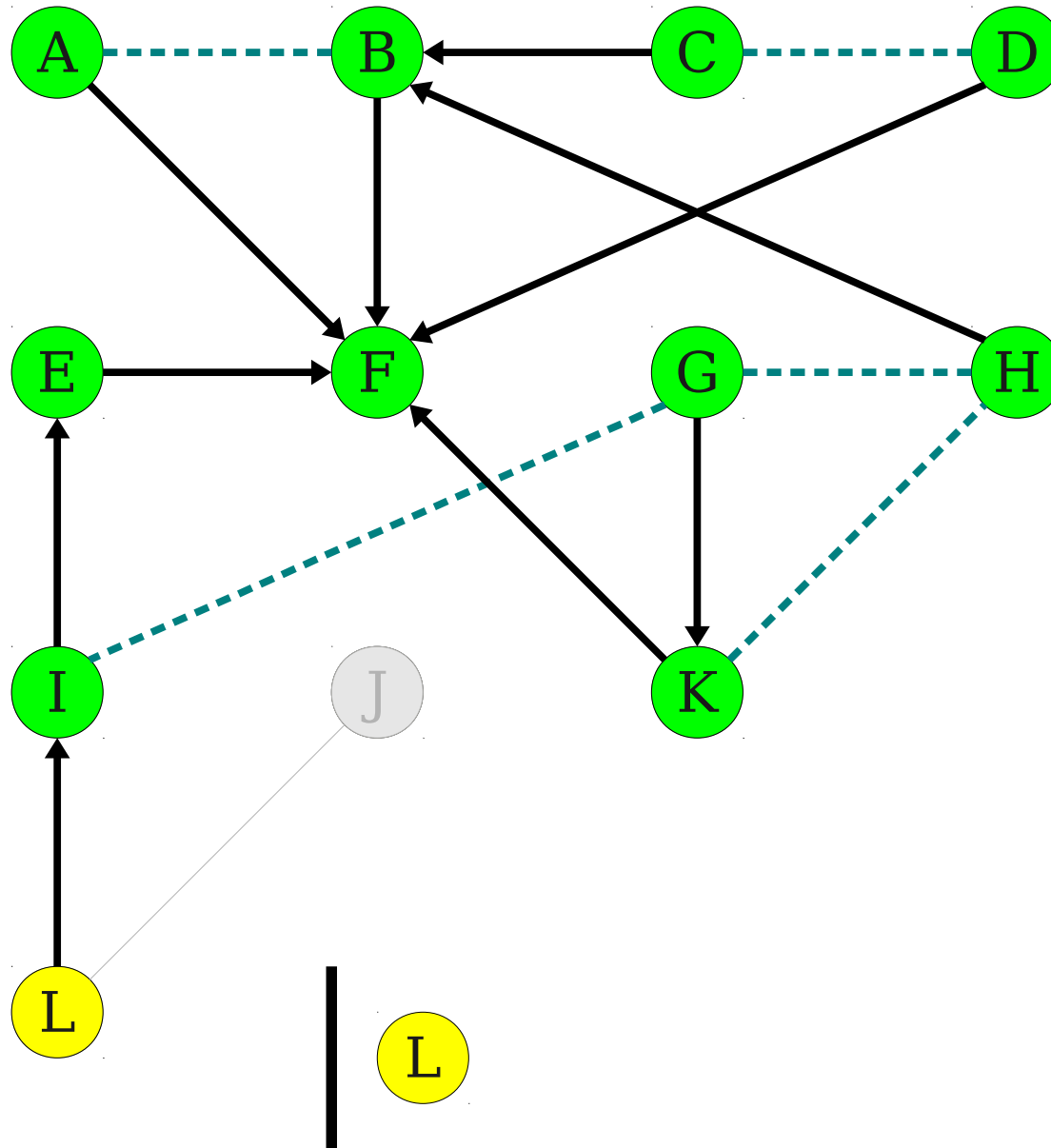
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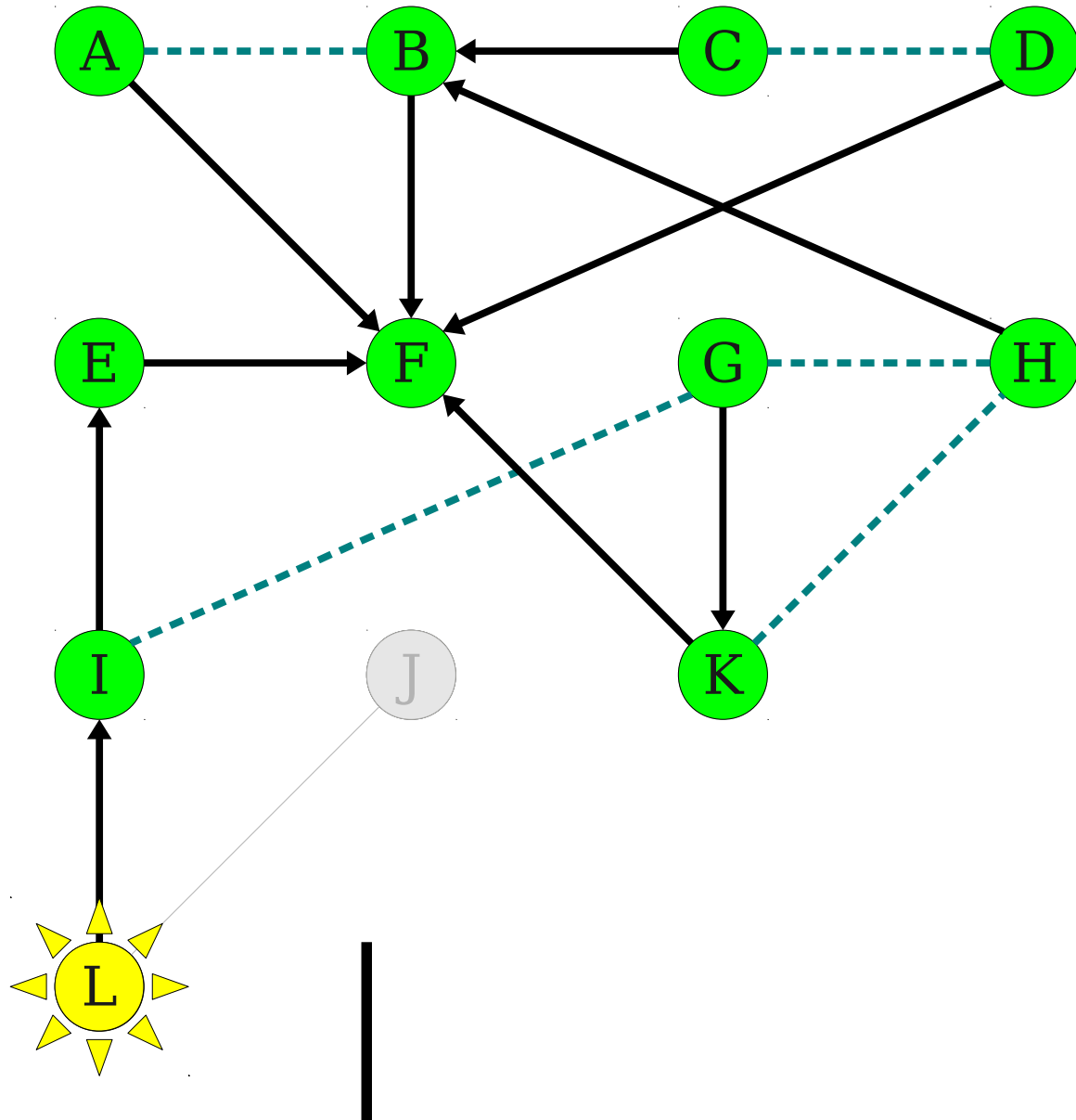
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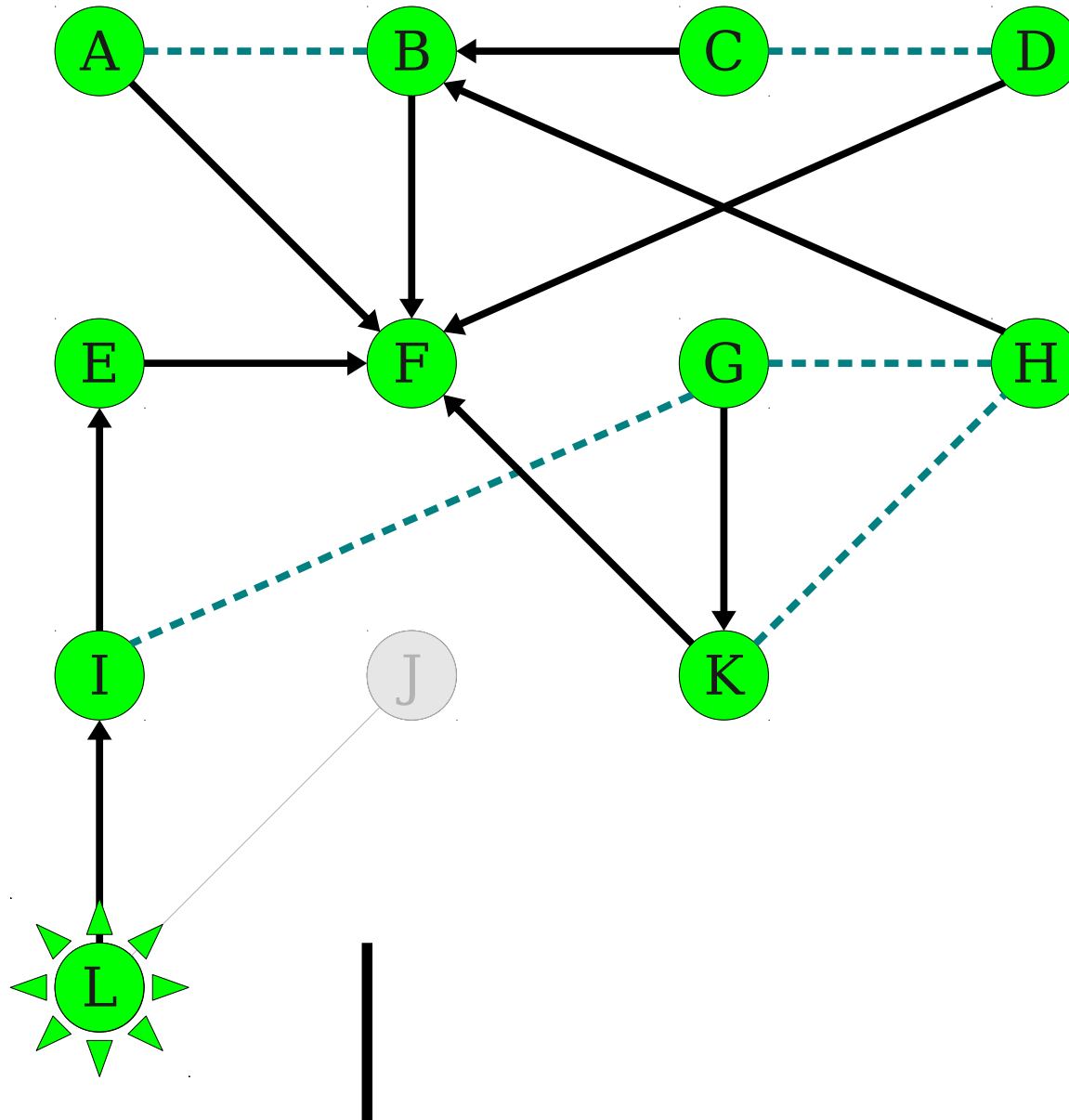
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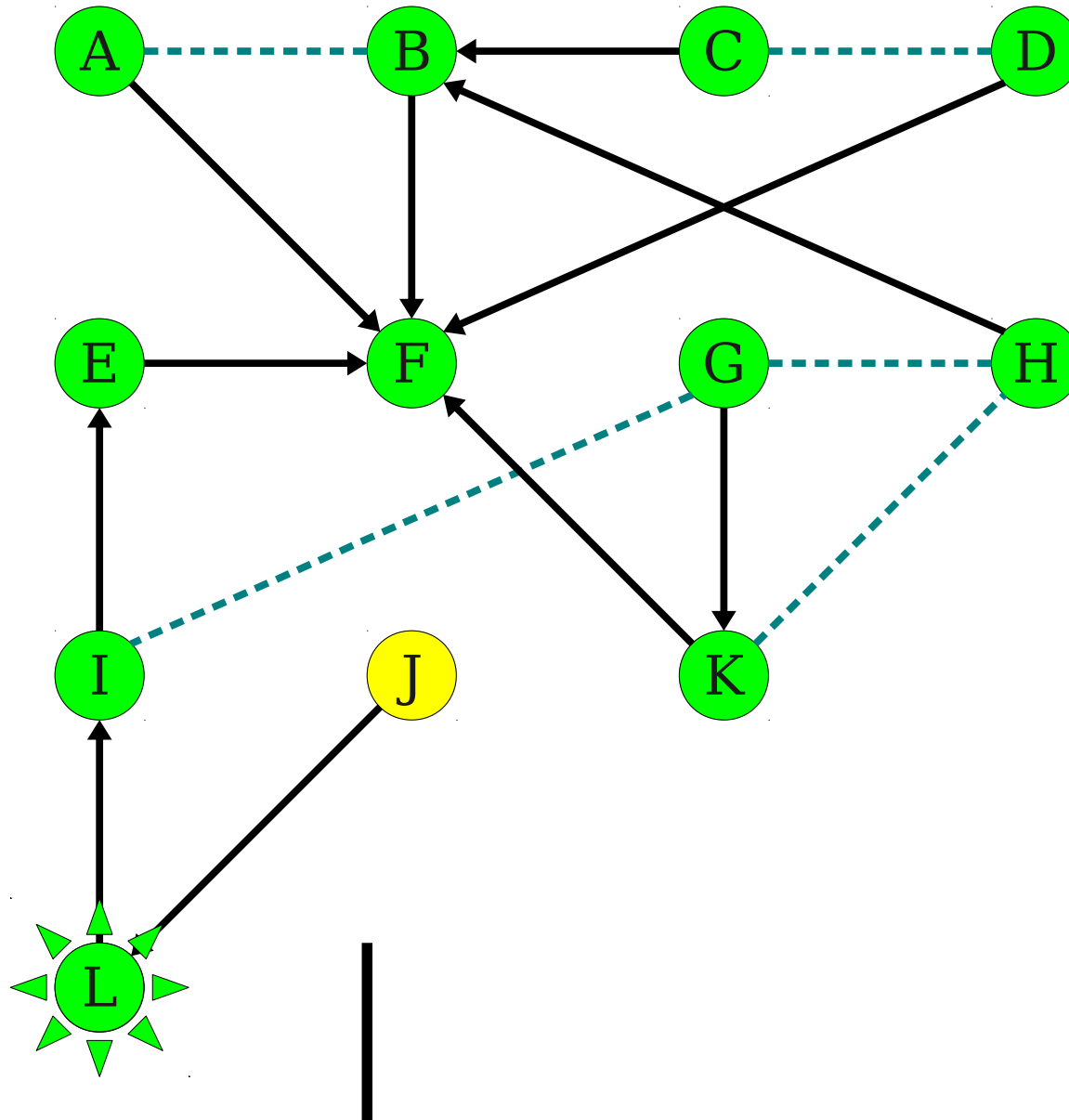
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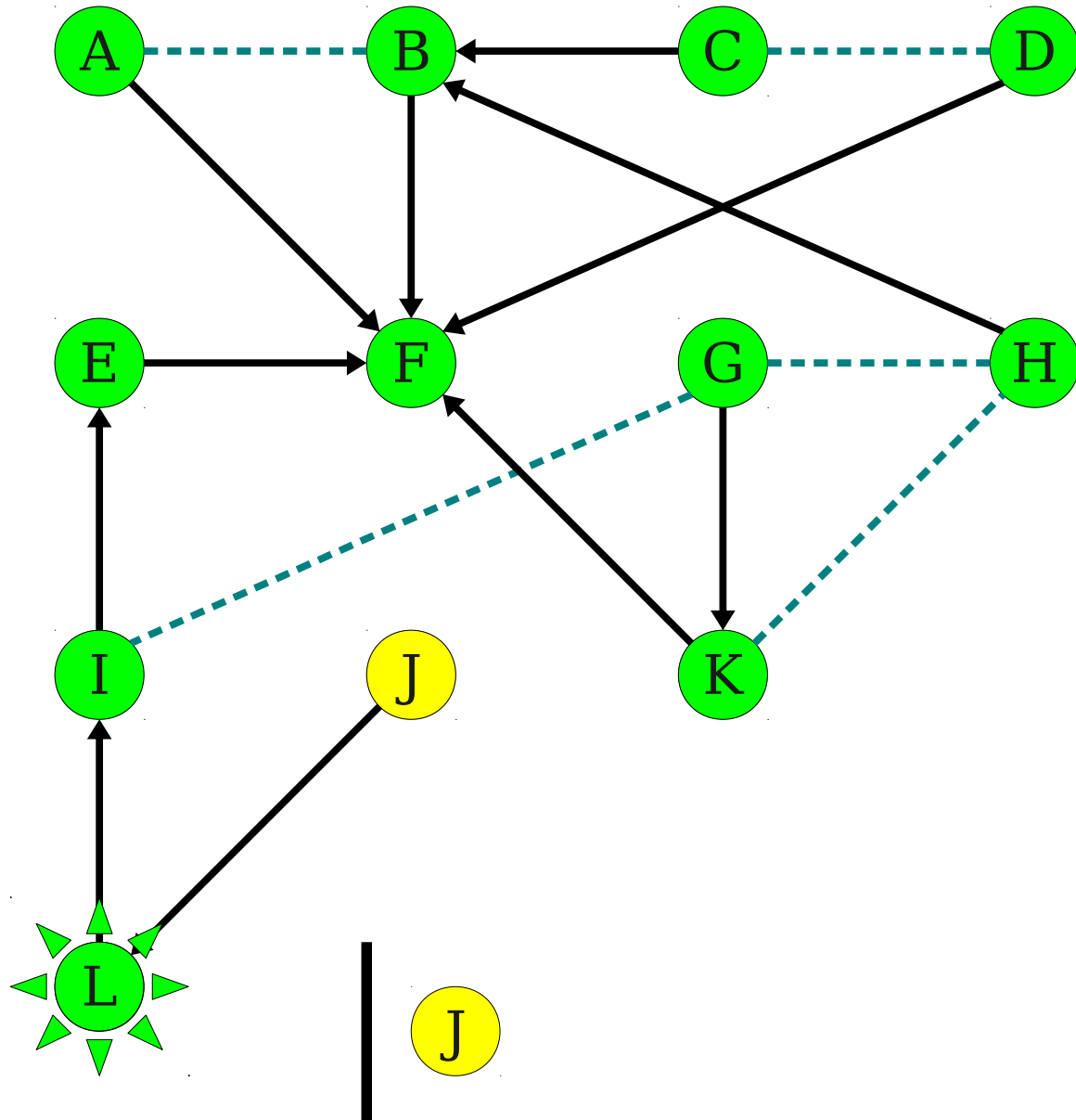
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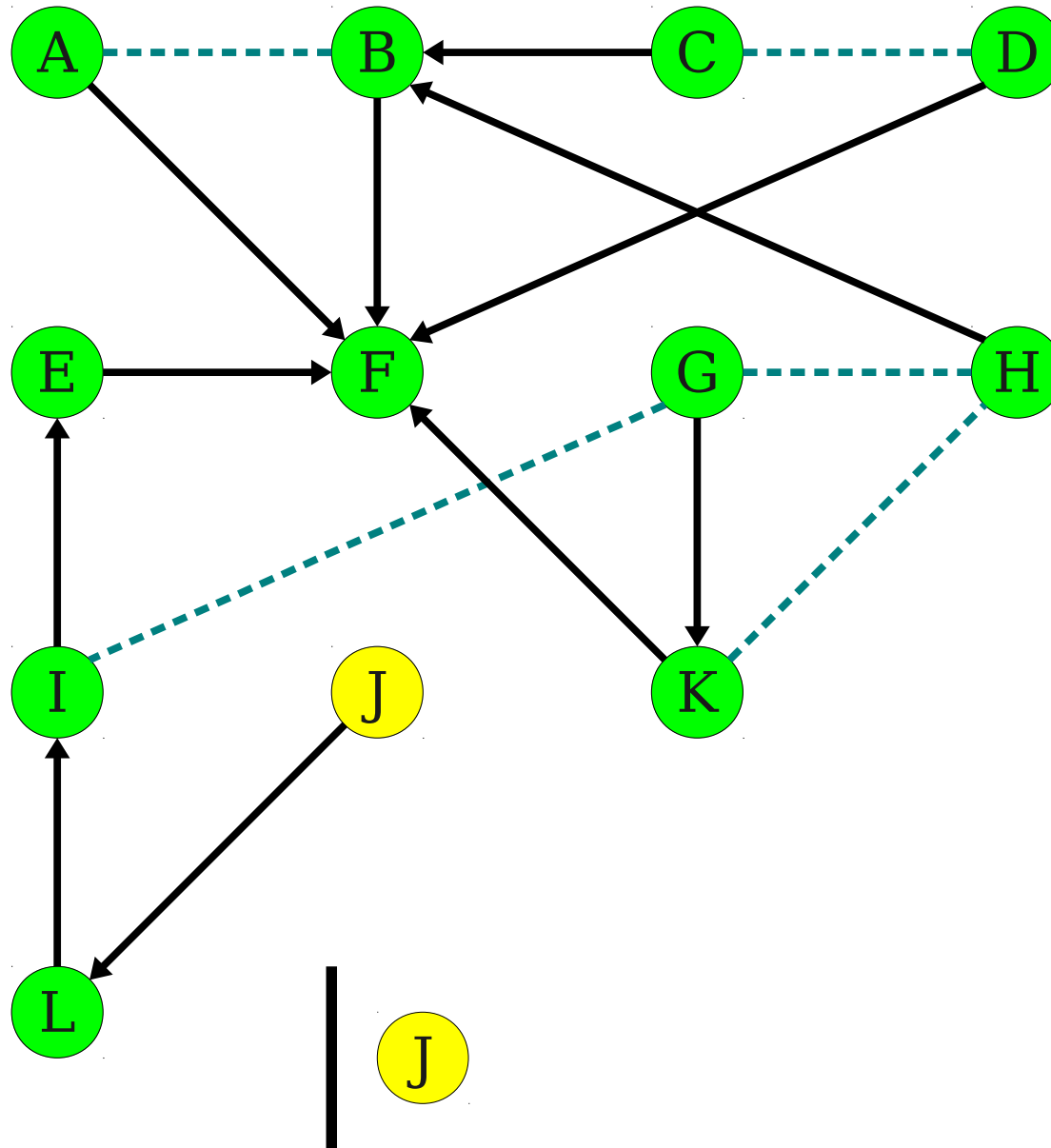
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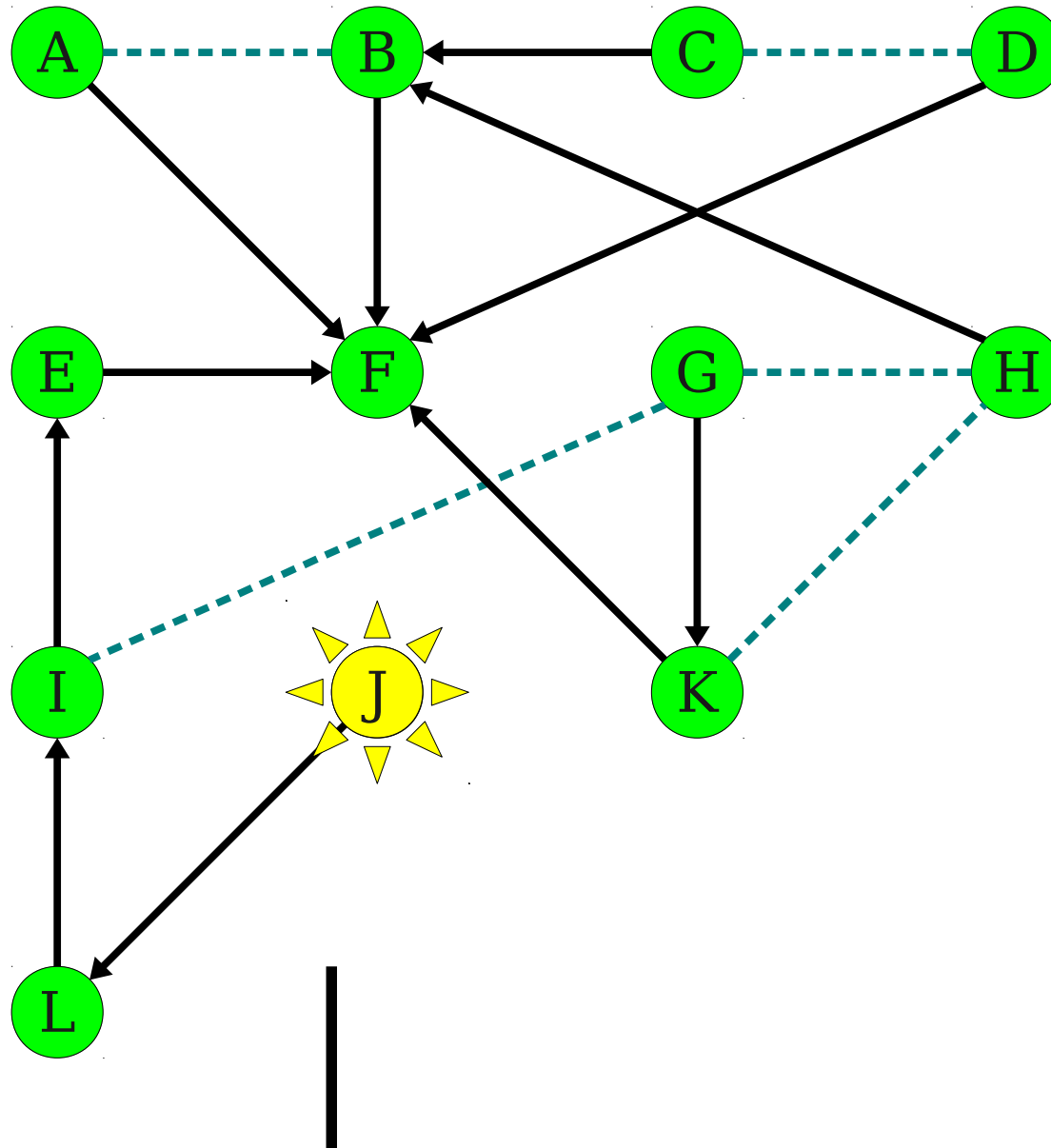
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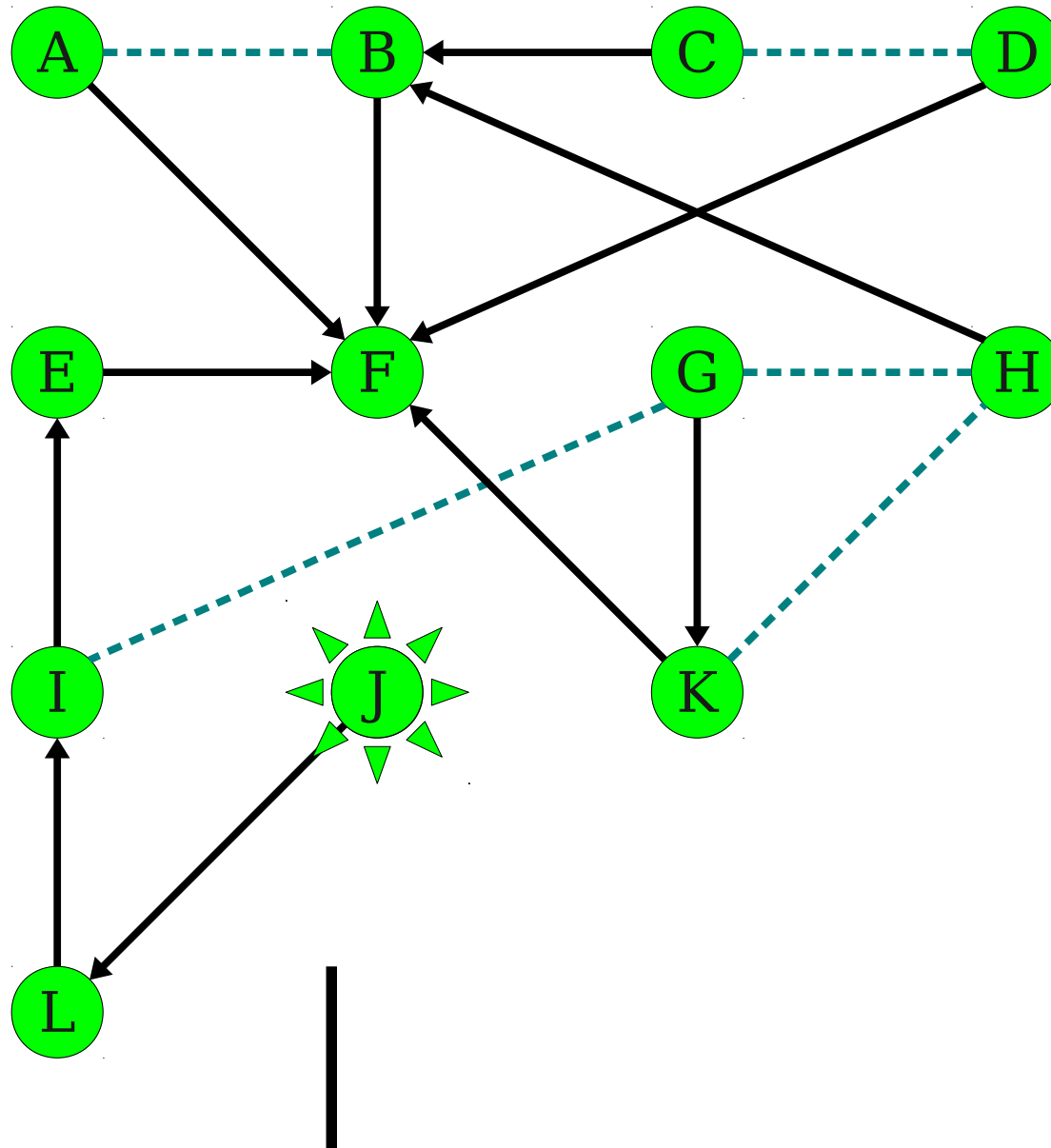
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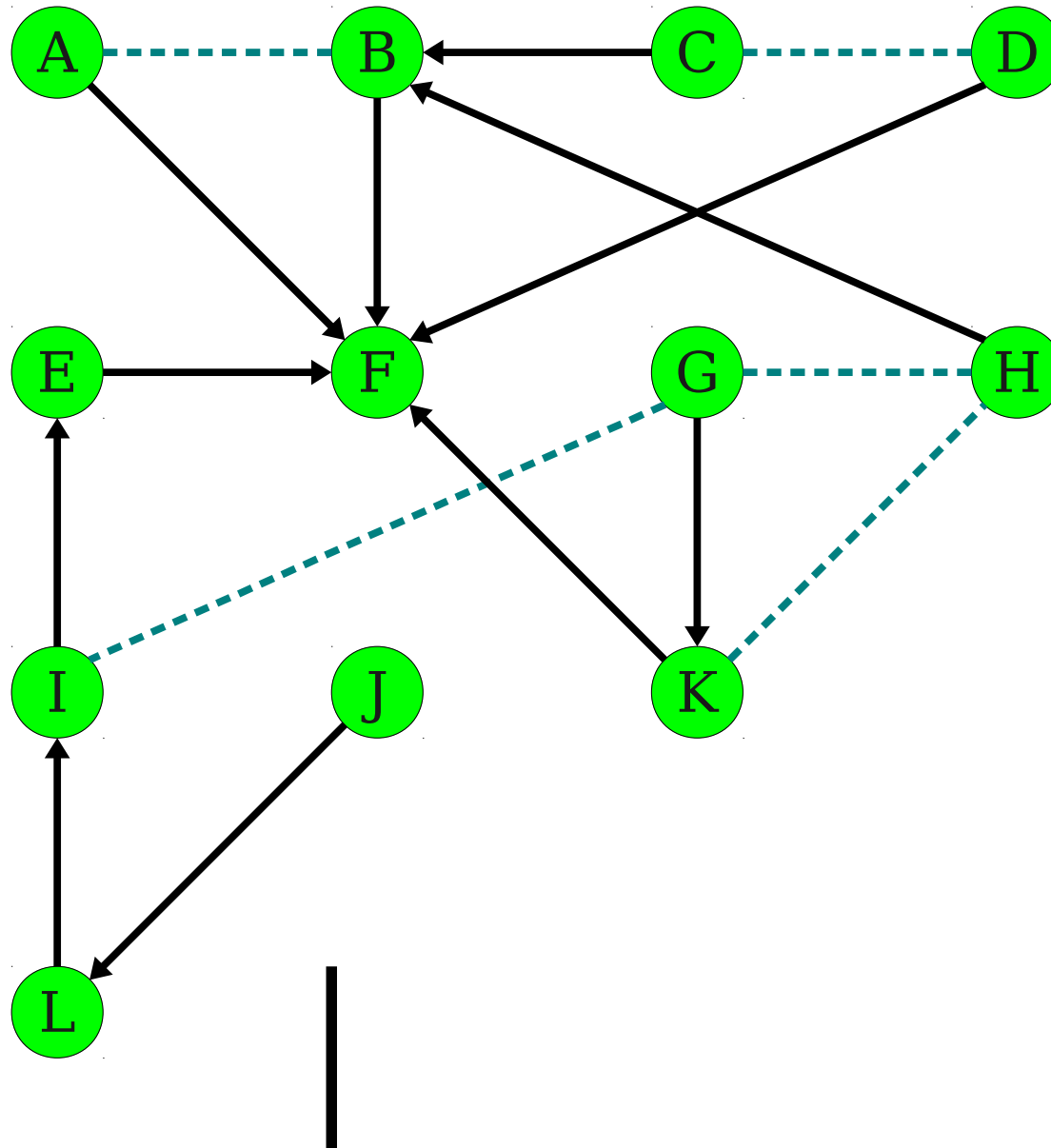
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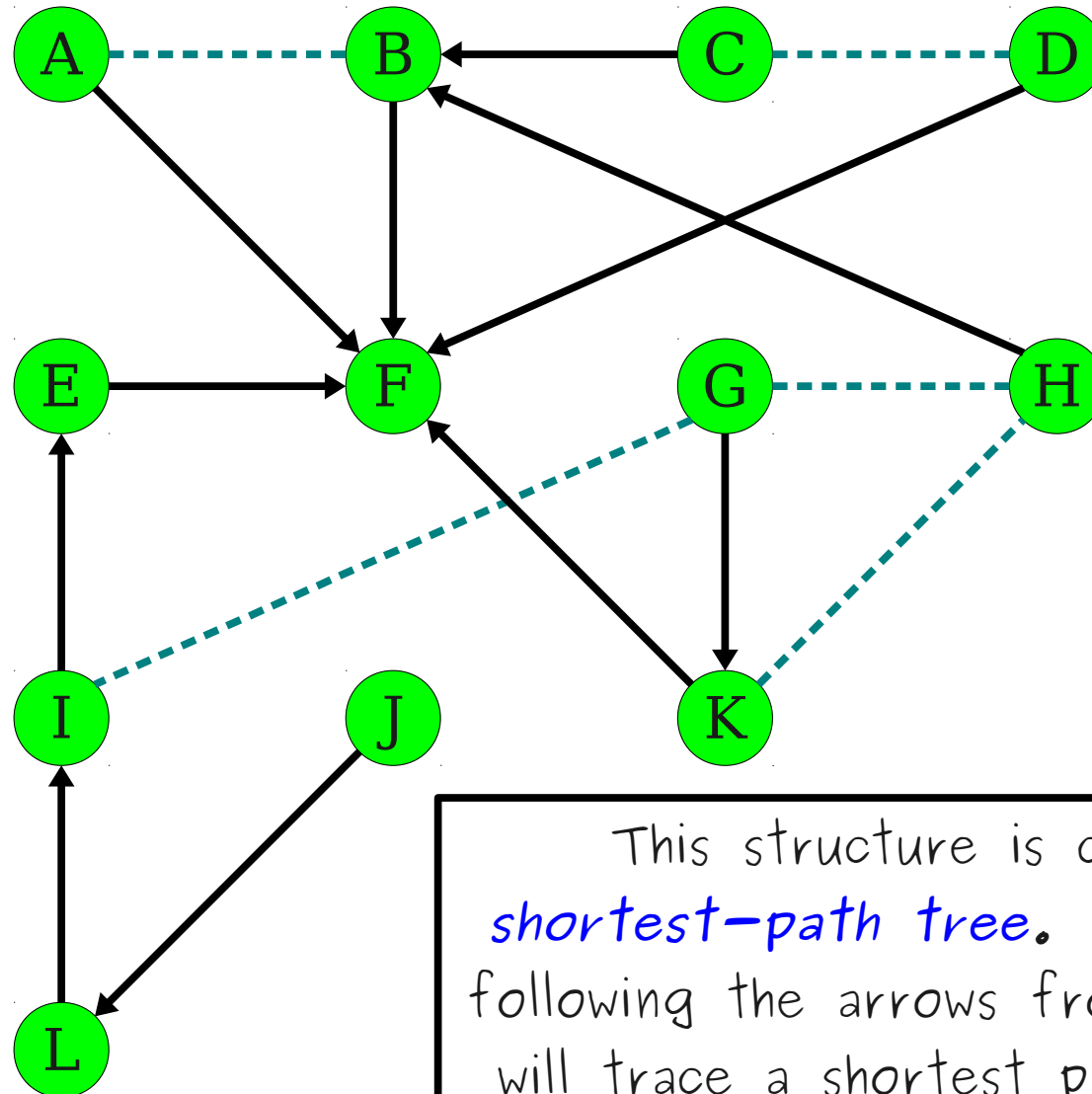
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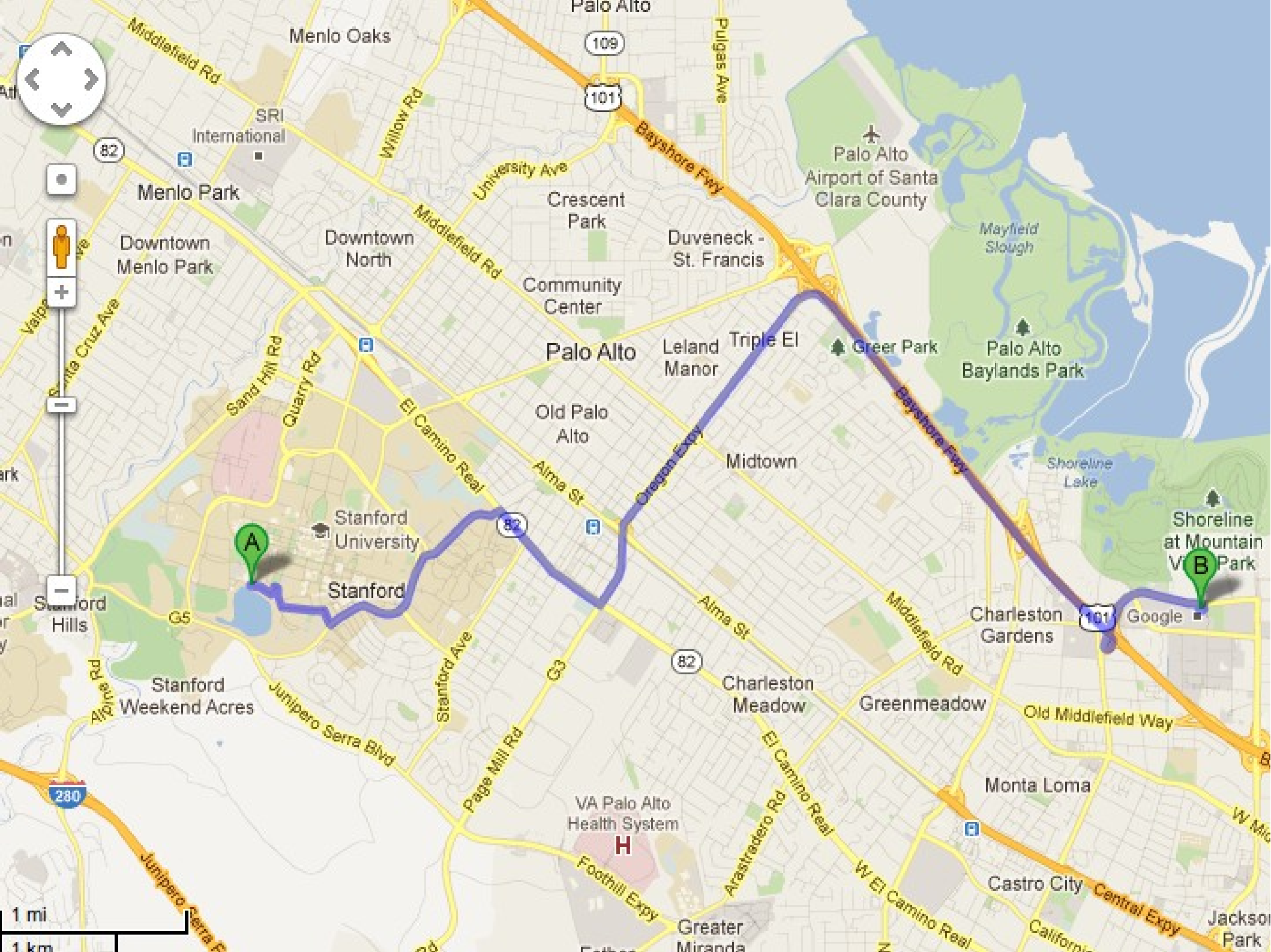
Breadth-First Search



Breadth-First Search



This structure is called a *shortest-path tree*. Notice how following the arrows from any node will trace a shortest path back to the root in reverse.

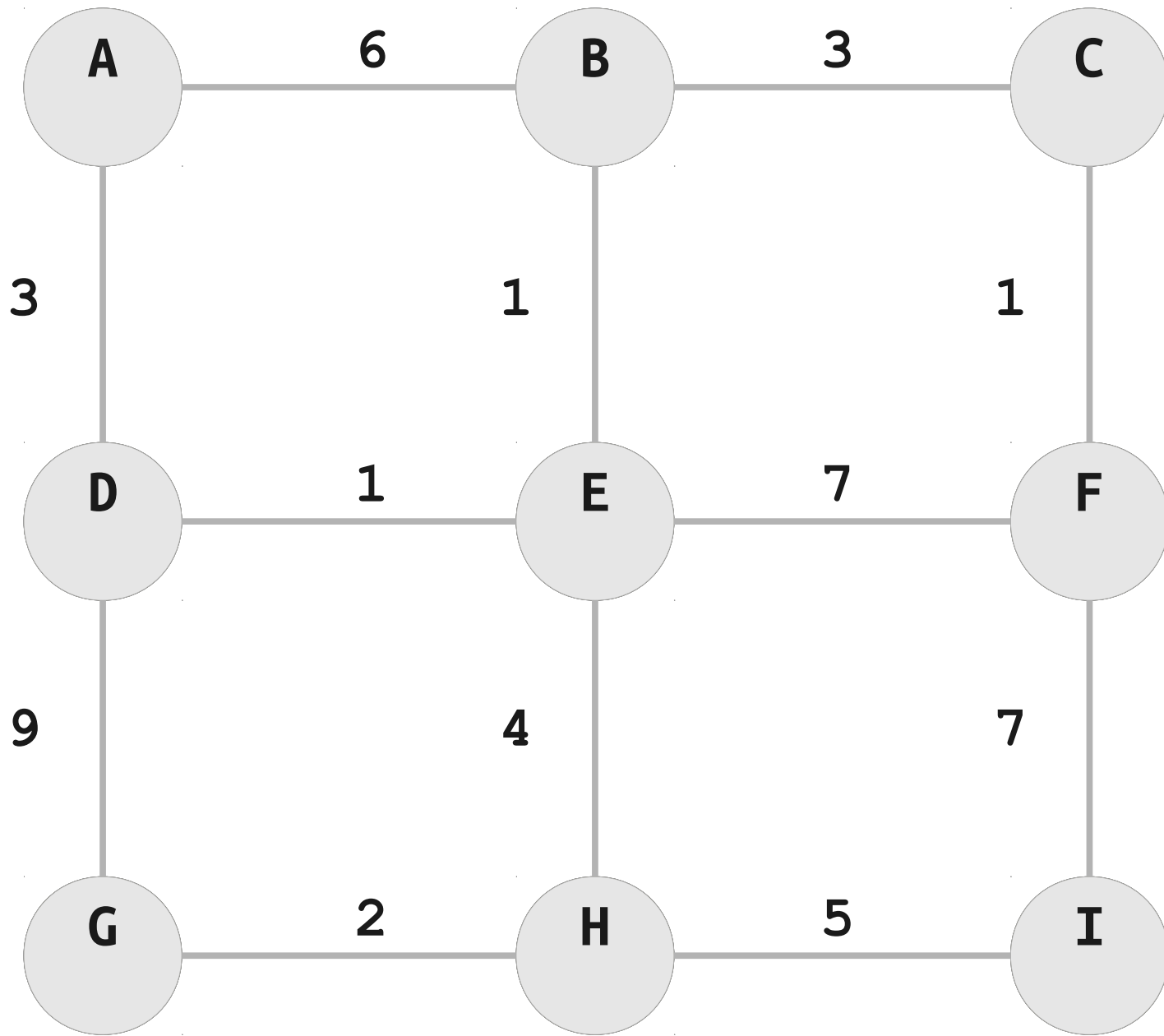


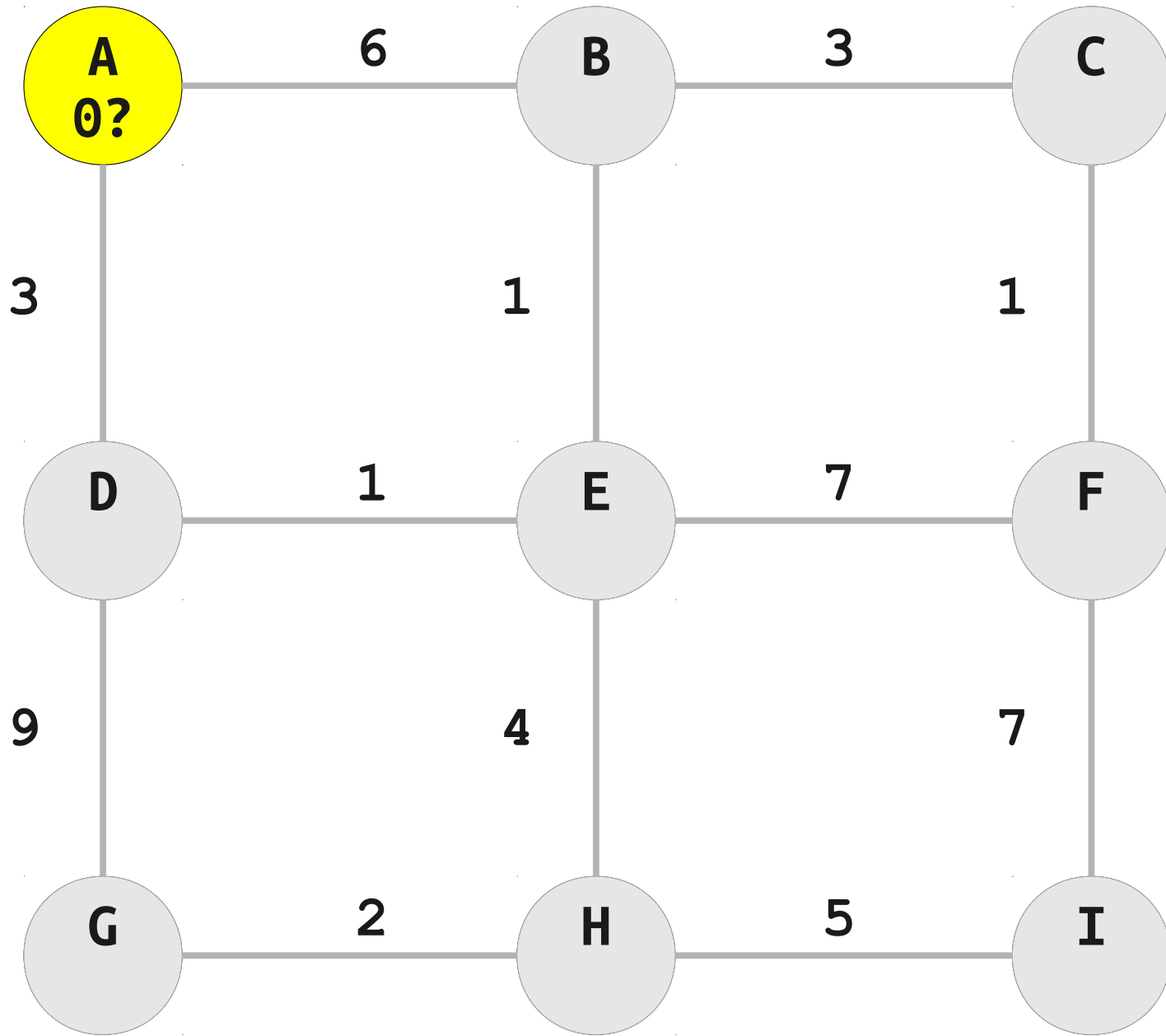
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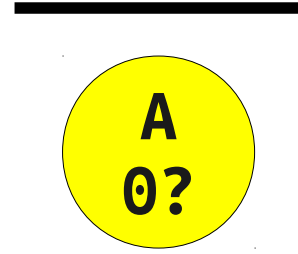
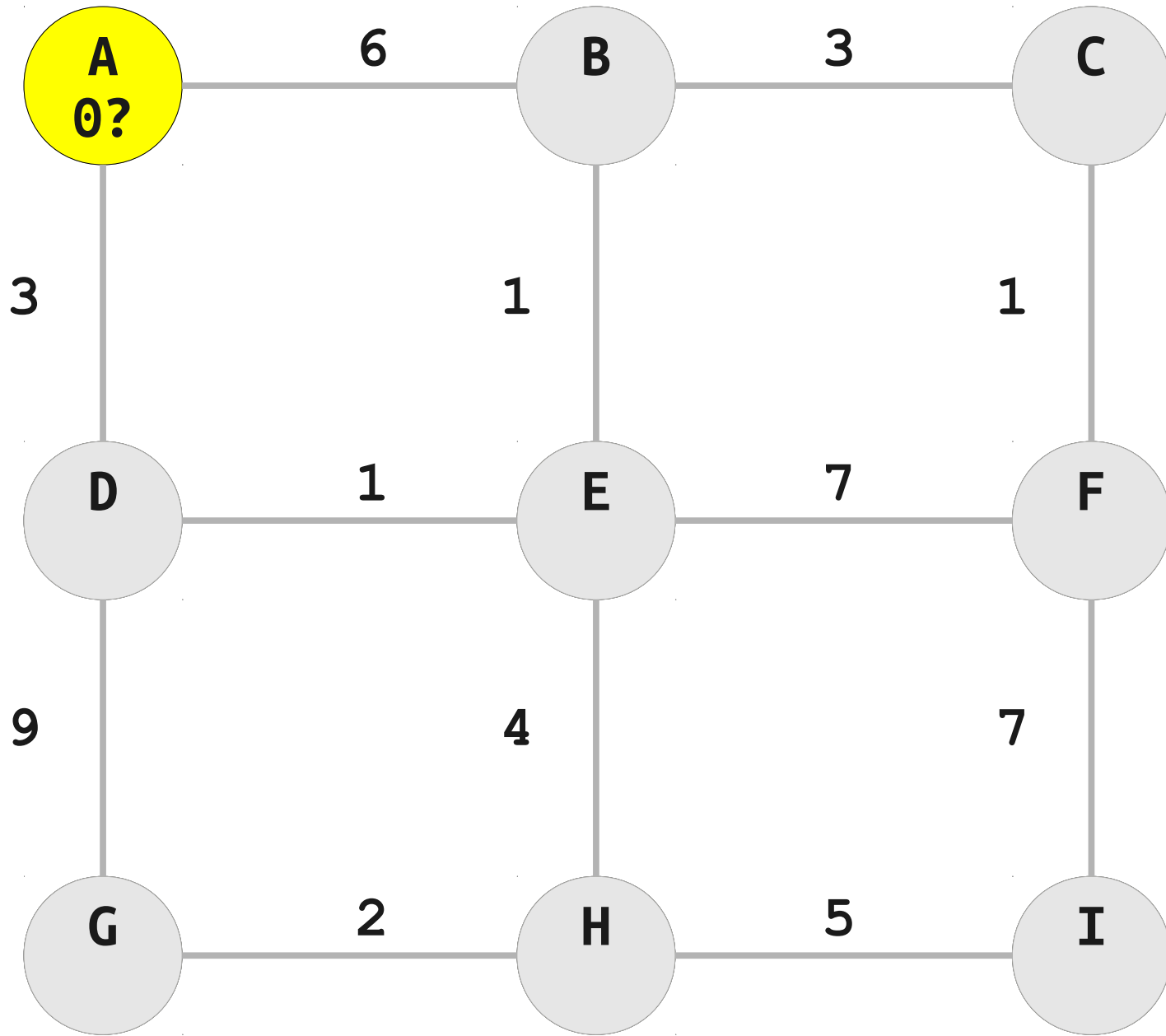
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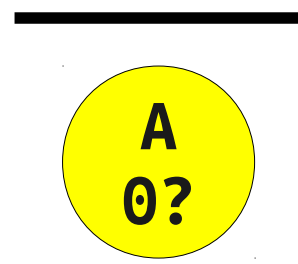
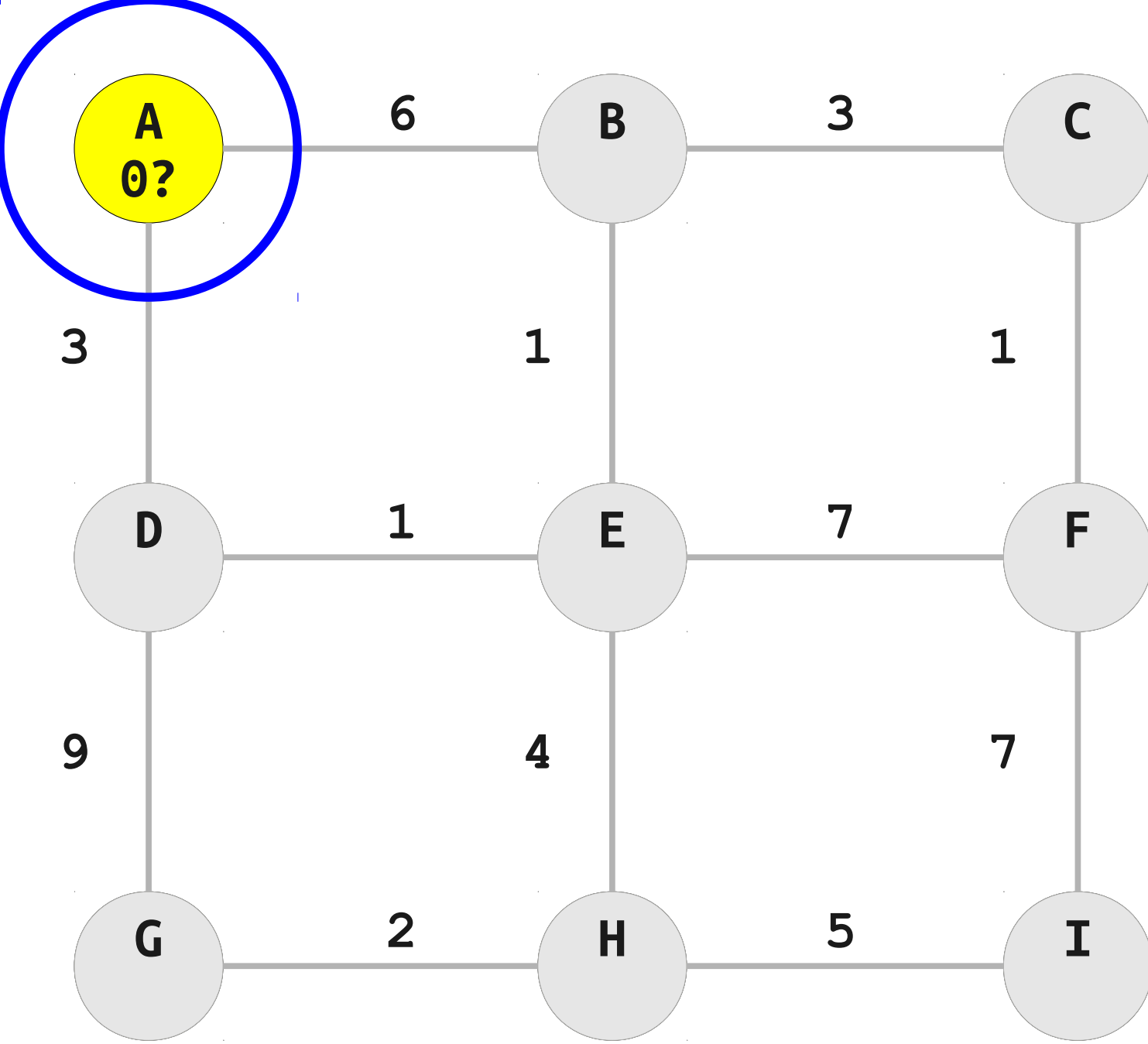
The Shortest Path Problem

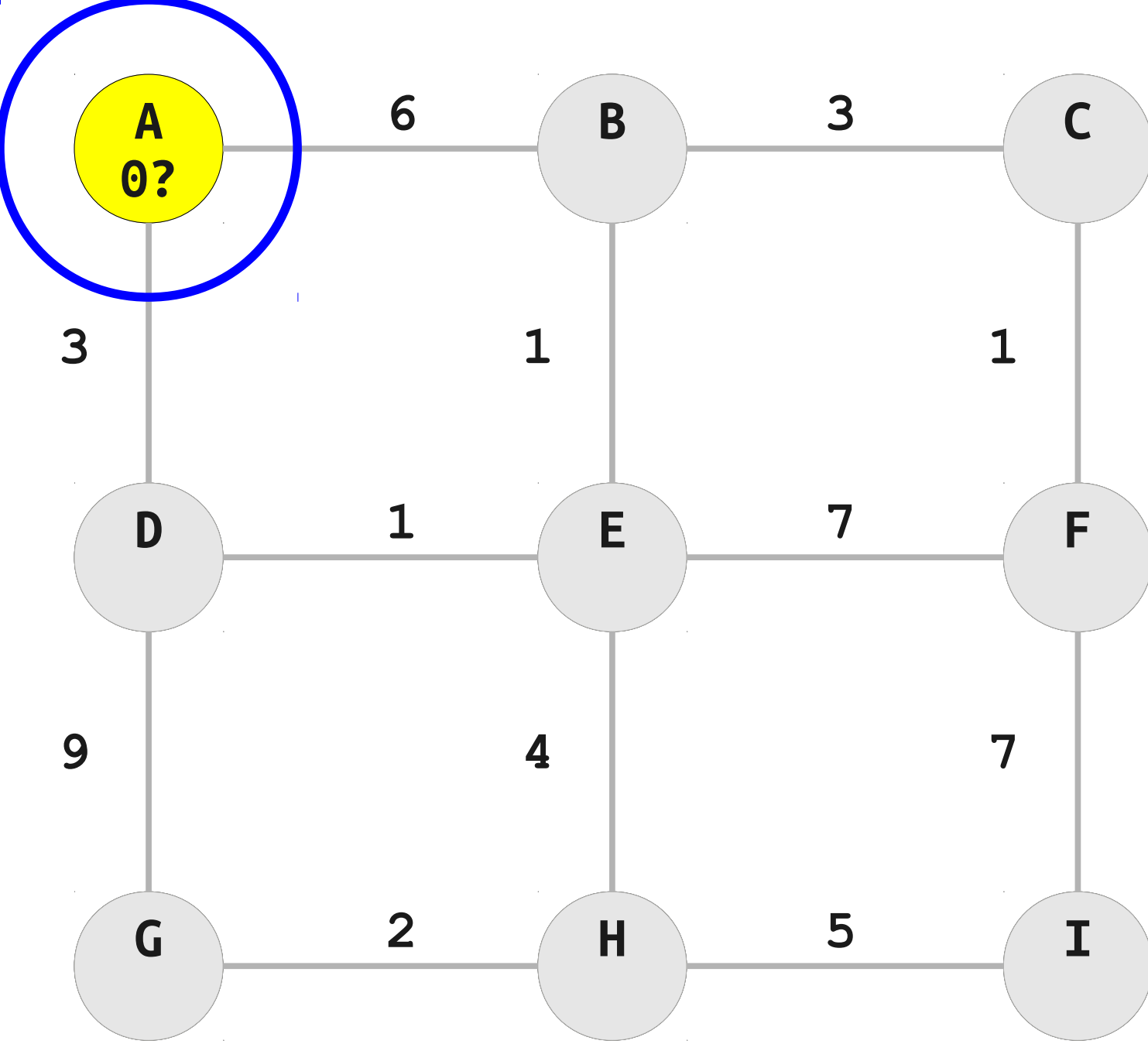
- Suppose that you have a graph representing different locations.
- Each edge has an associated *cost* (or *weight*, *length*, etc.). We'll assume costs are nonnegative.
- Goal: Find the least-cost (or lowest-weight, lowest-length, etc.) path from some node u to a node v .

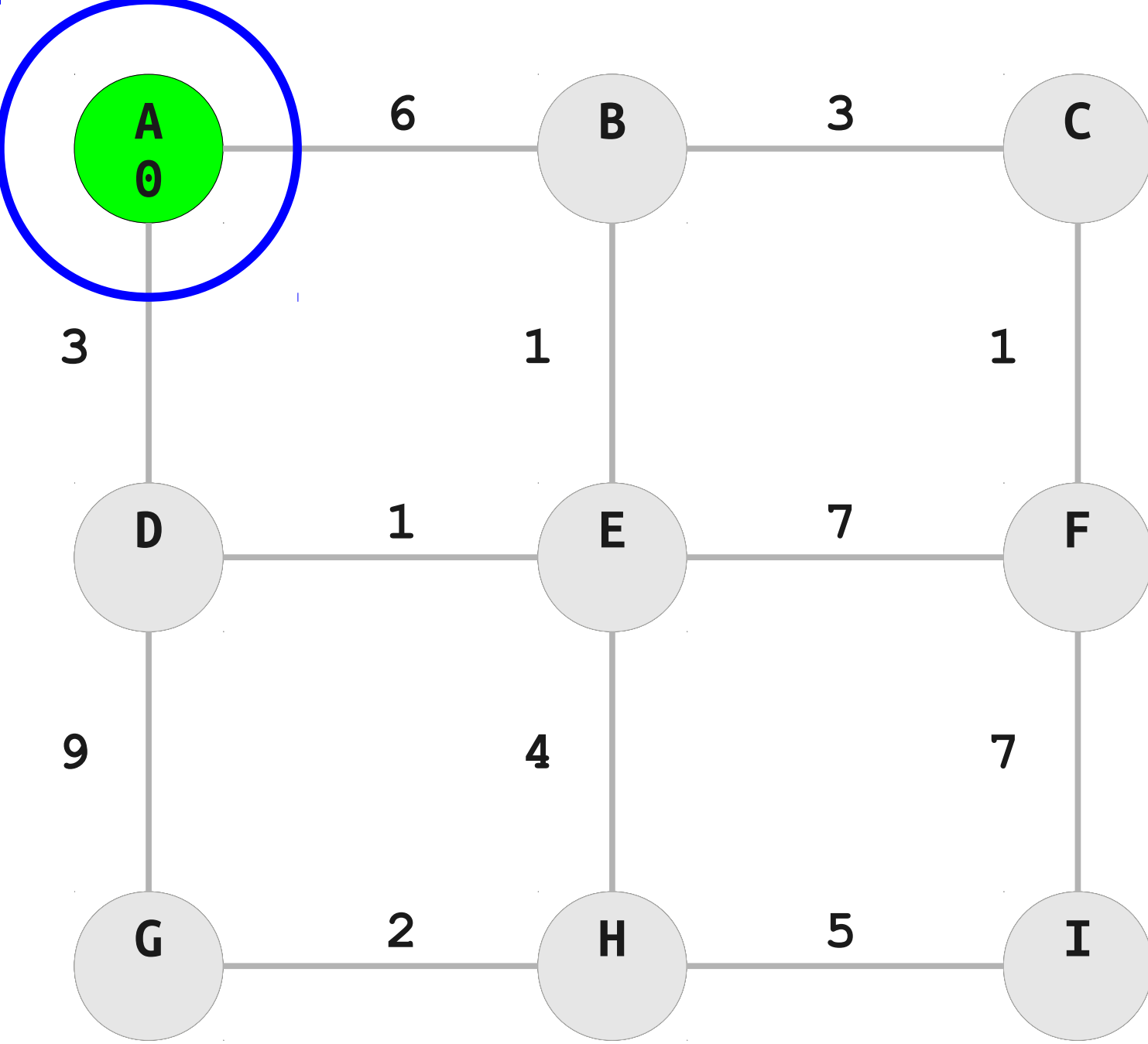


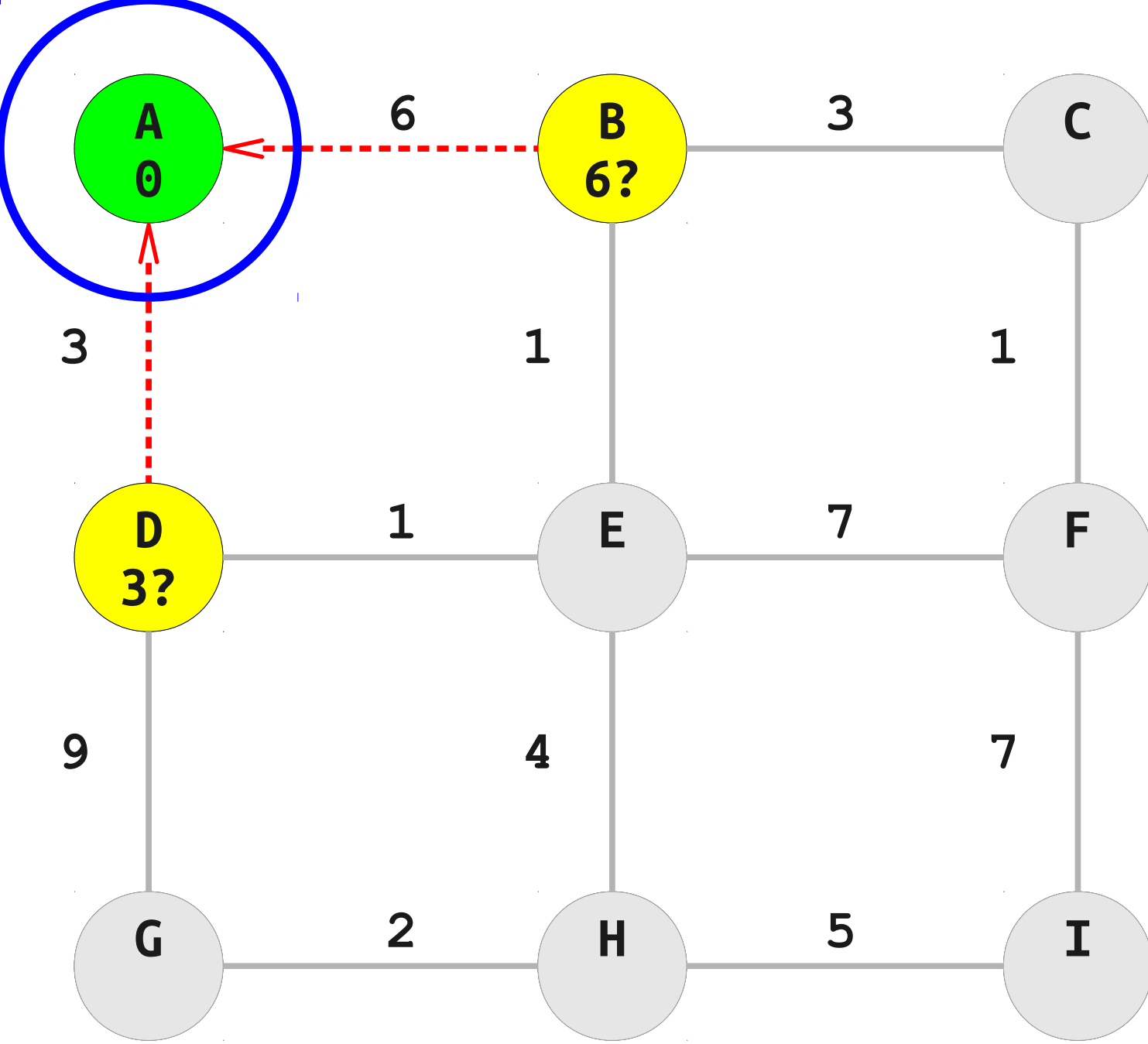


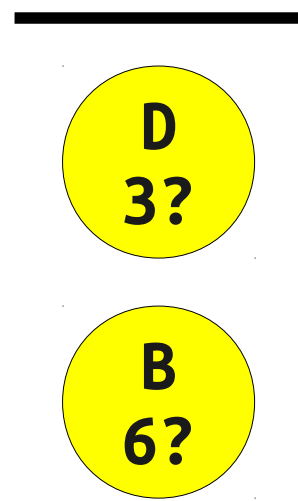
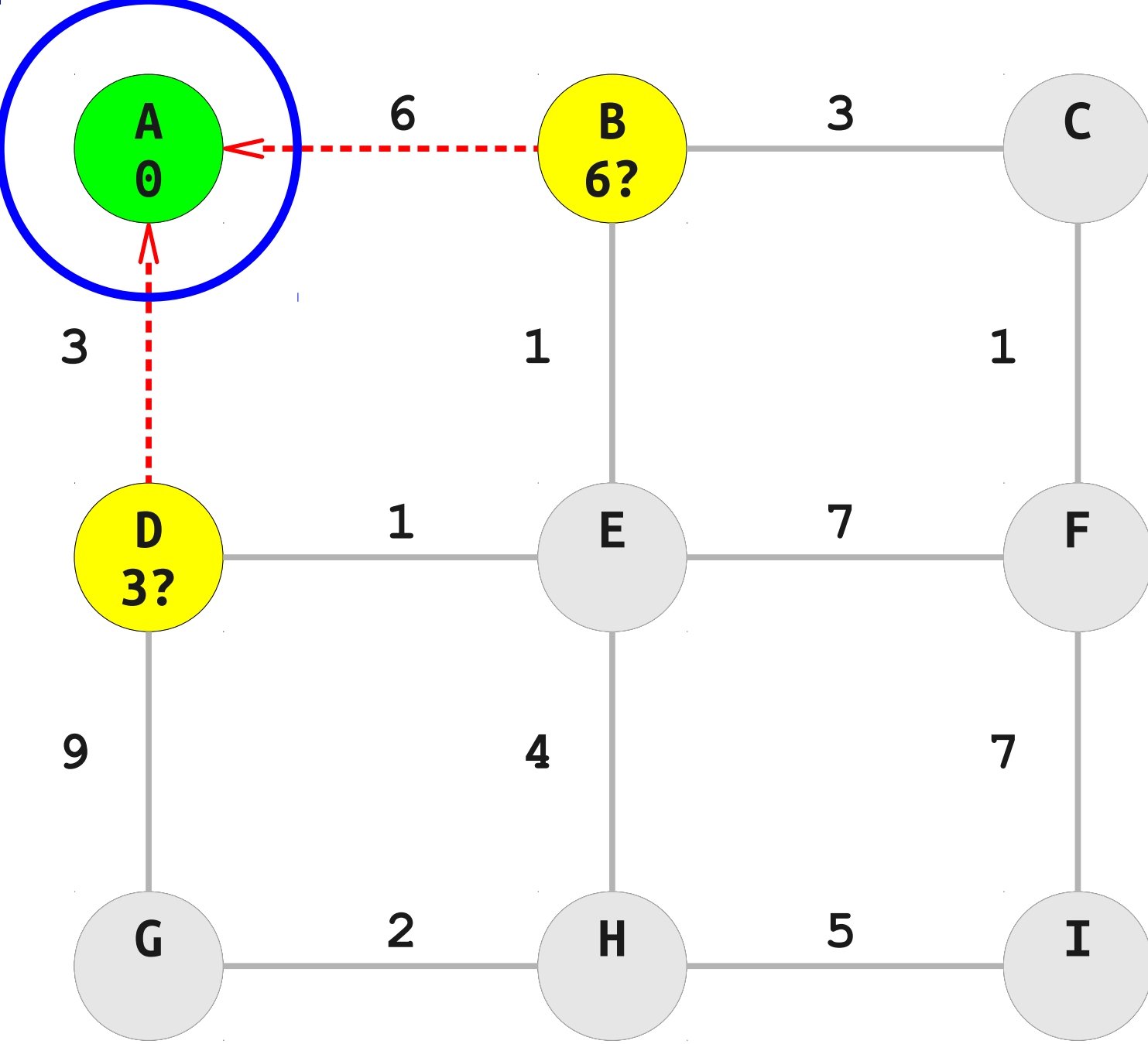


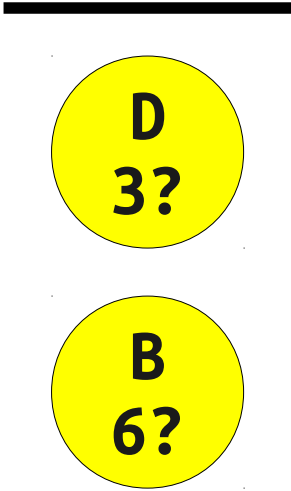
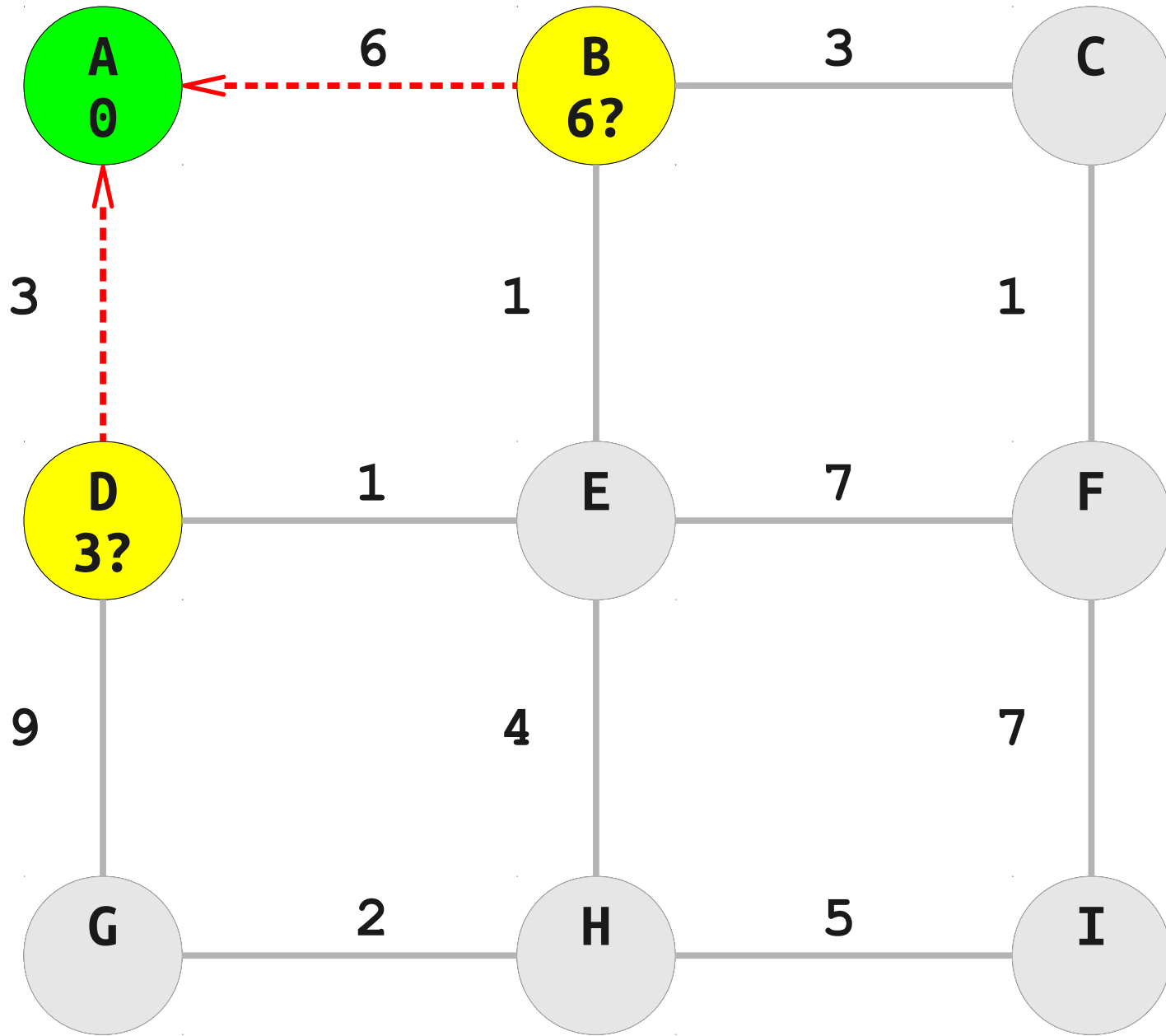


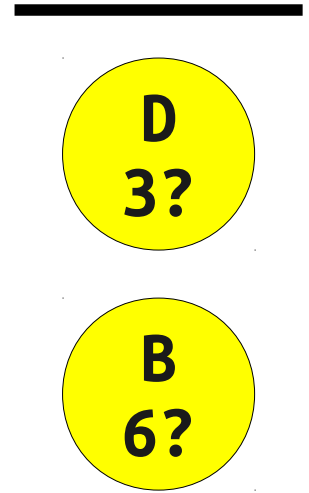
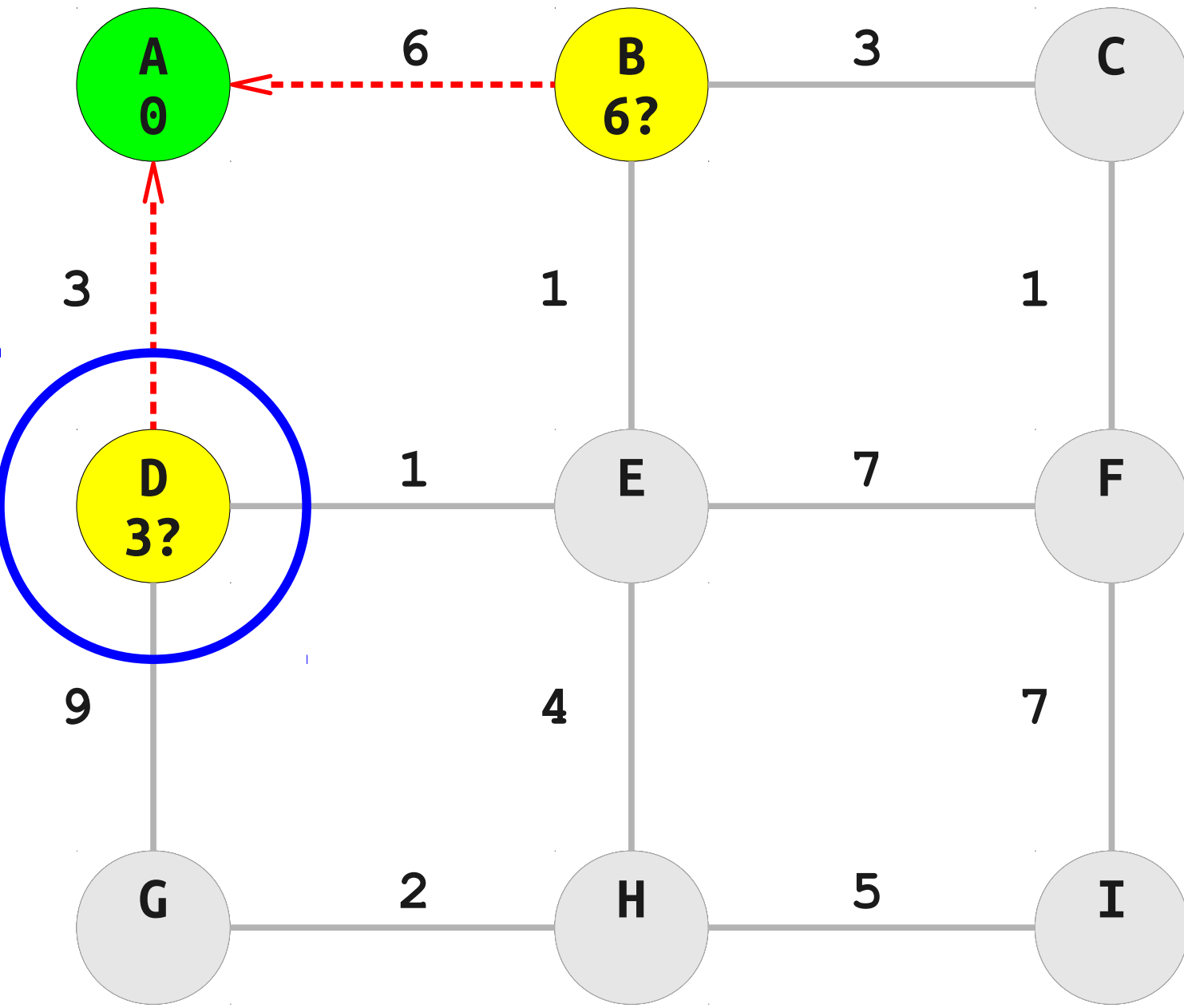


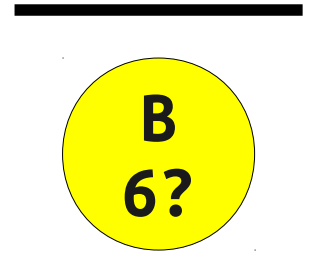
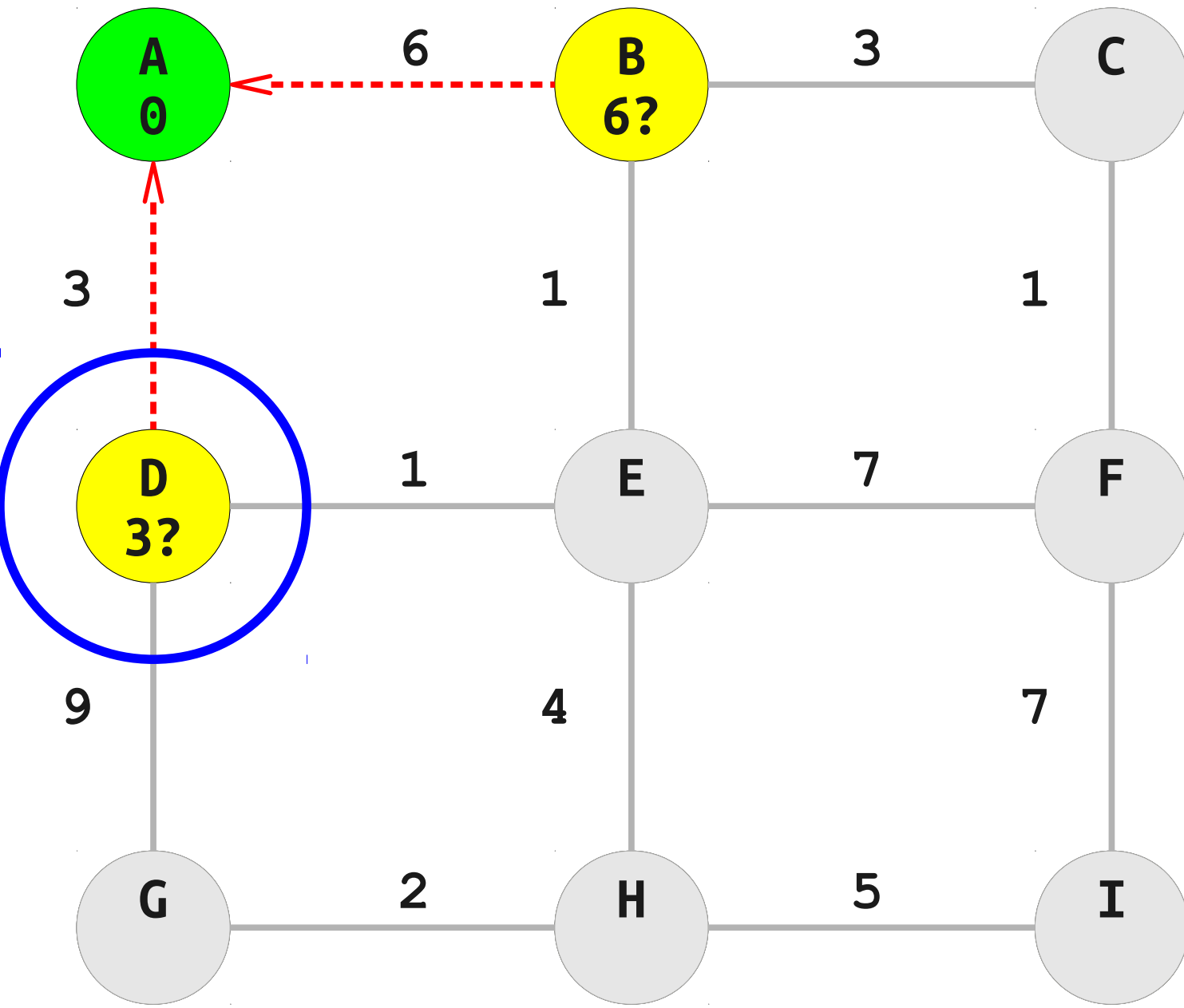


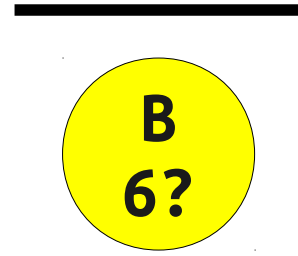
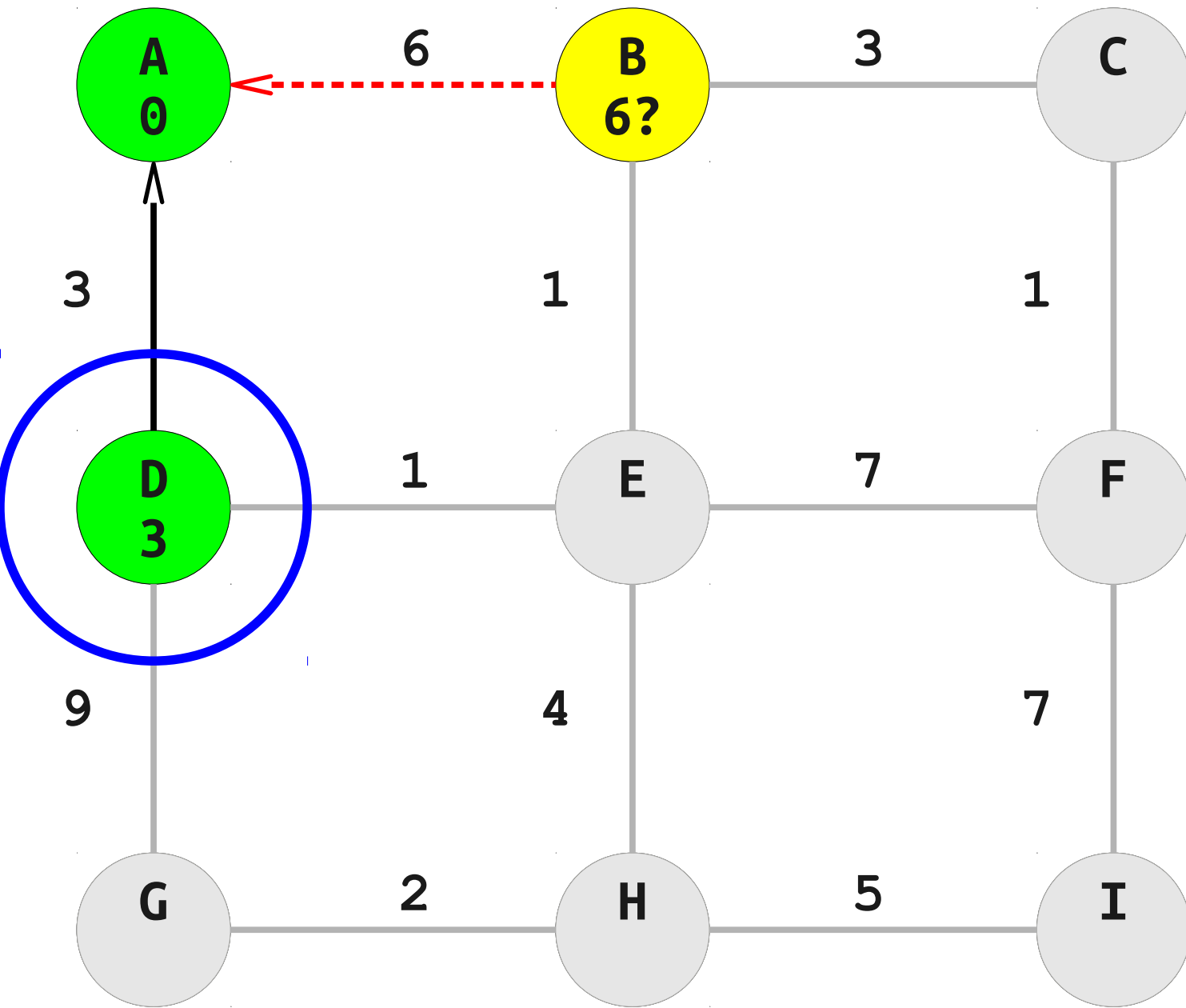


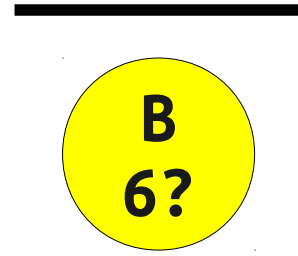
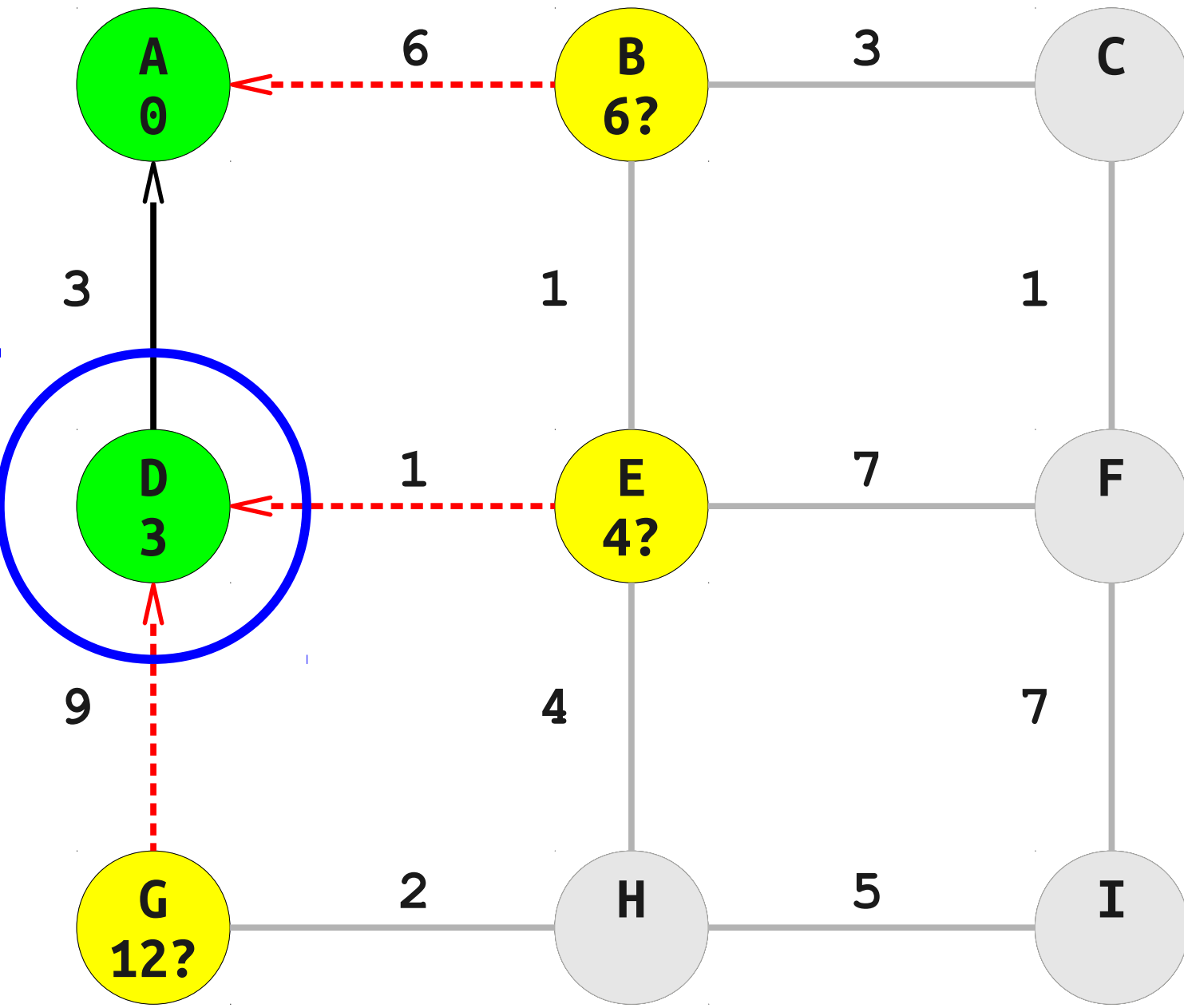


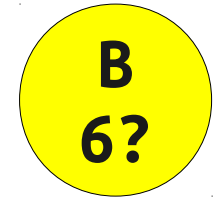
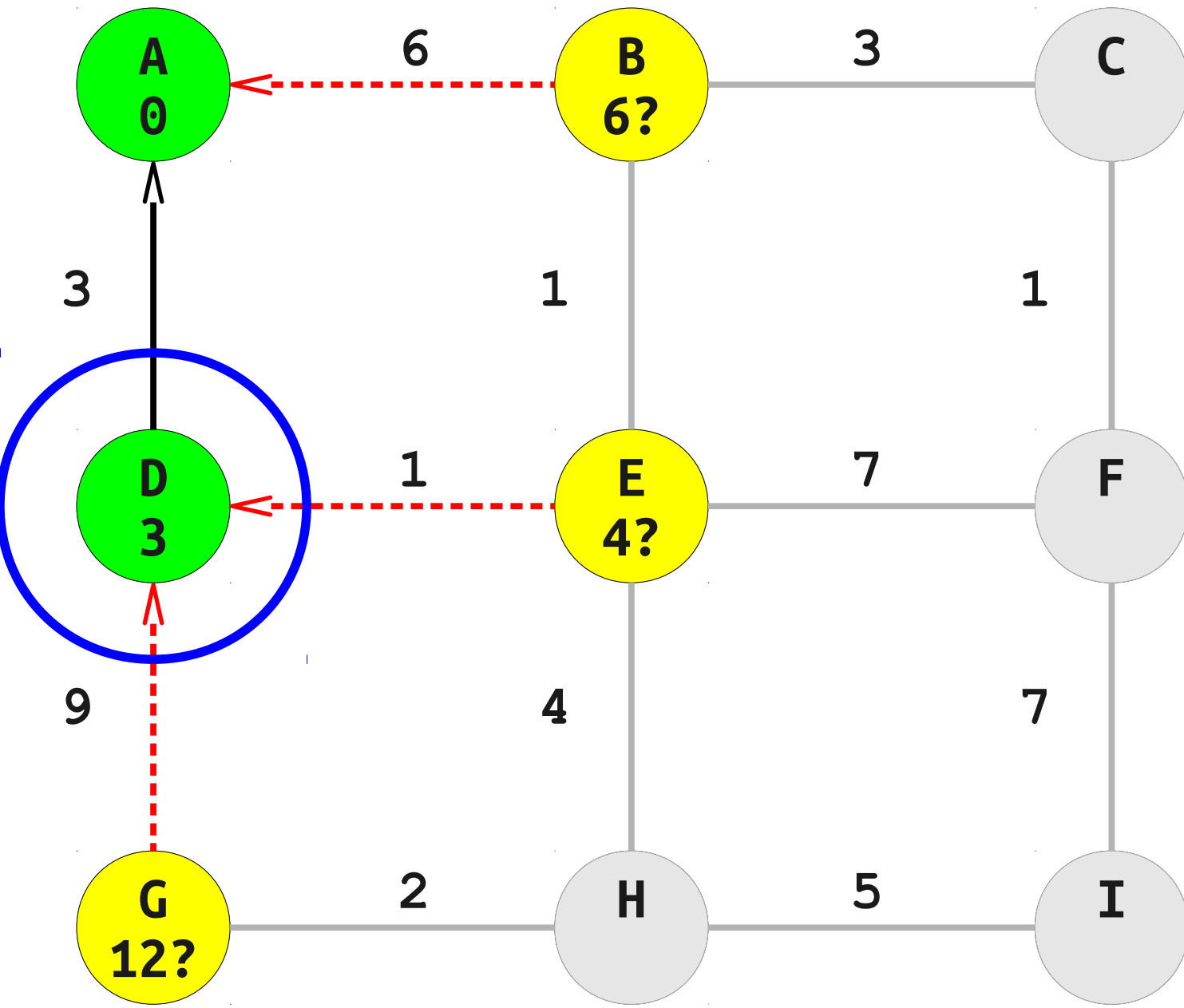


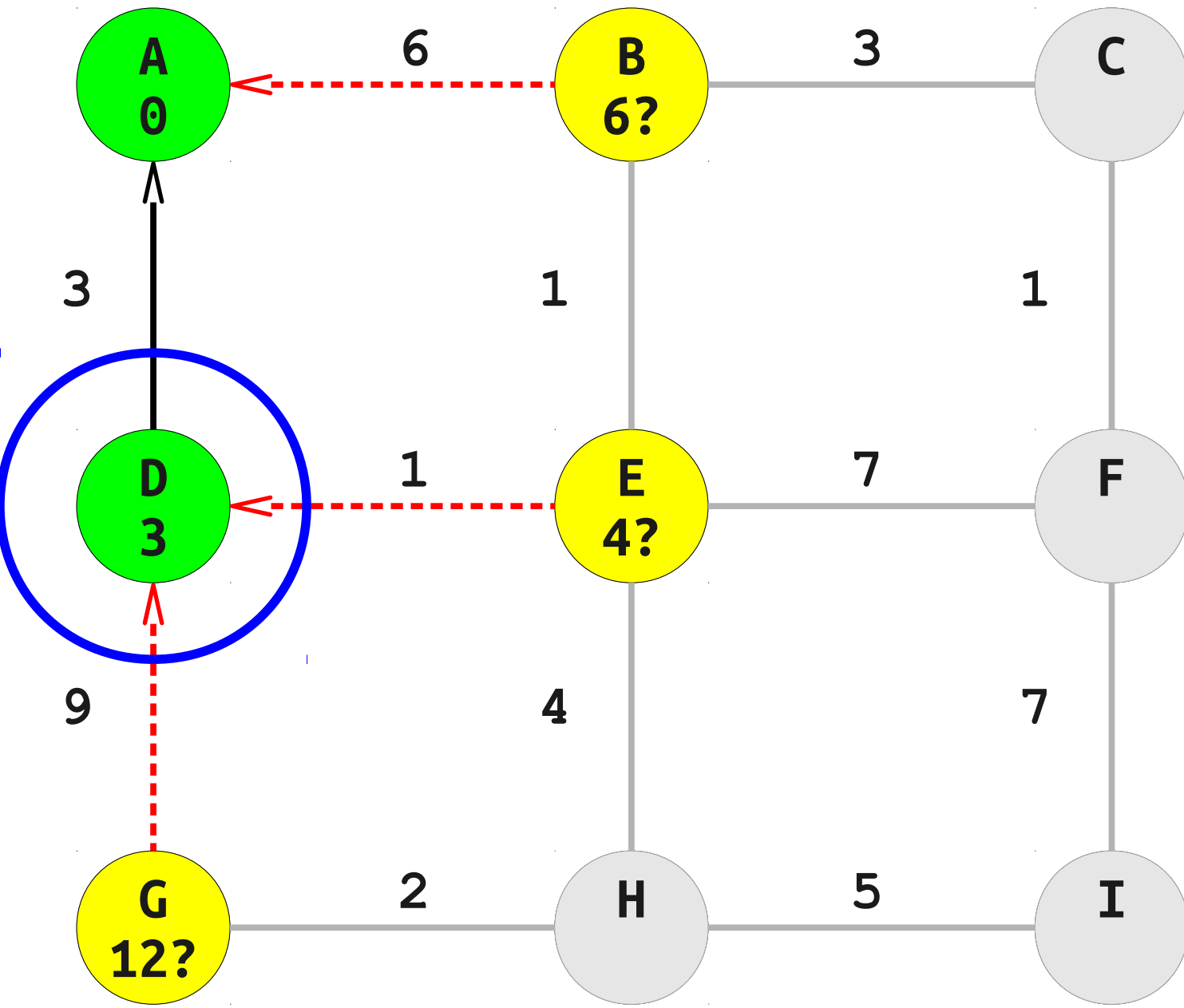








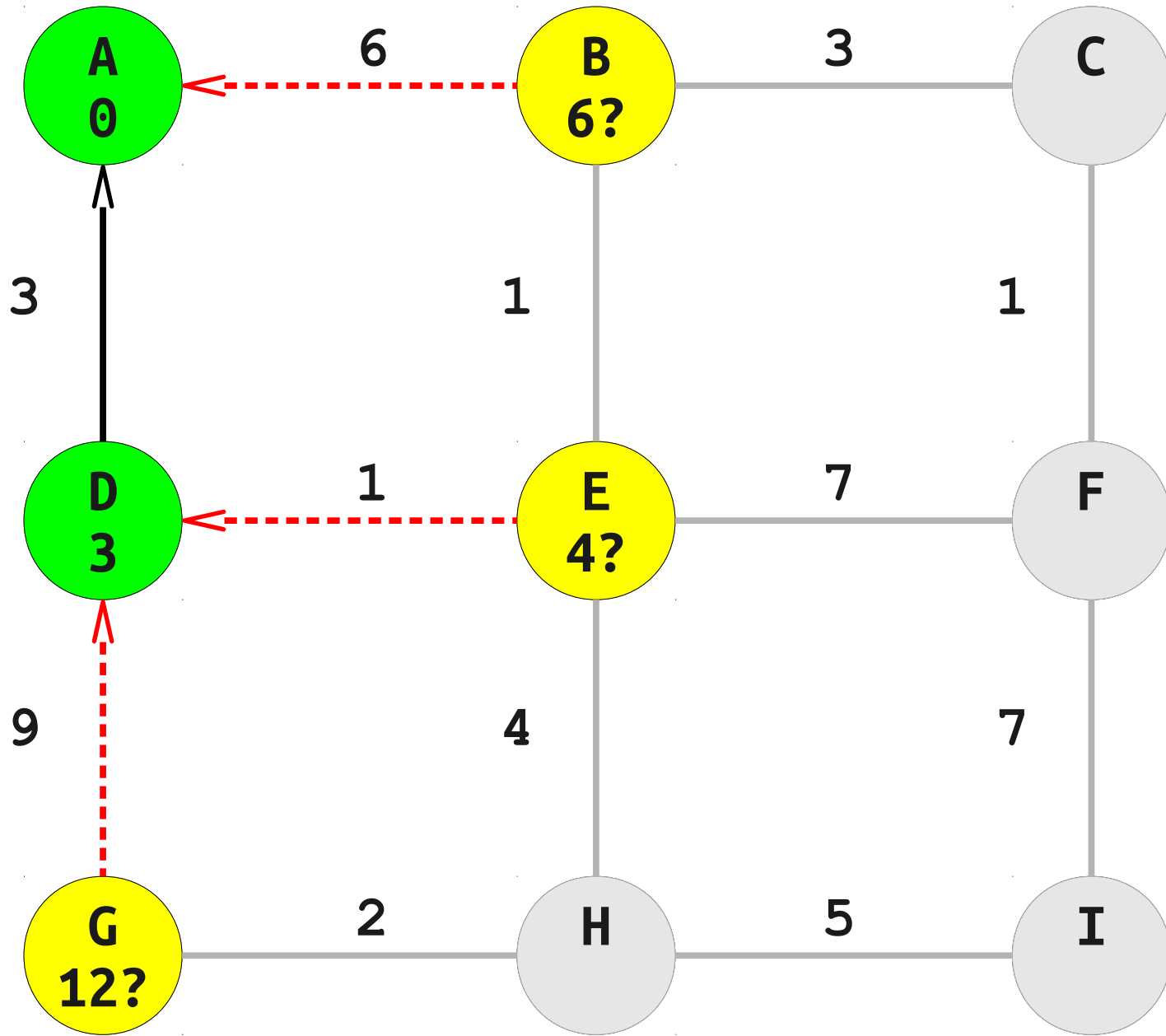




E
4?

B
6?

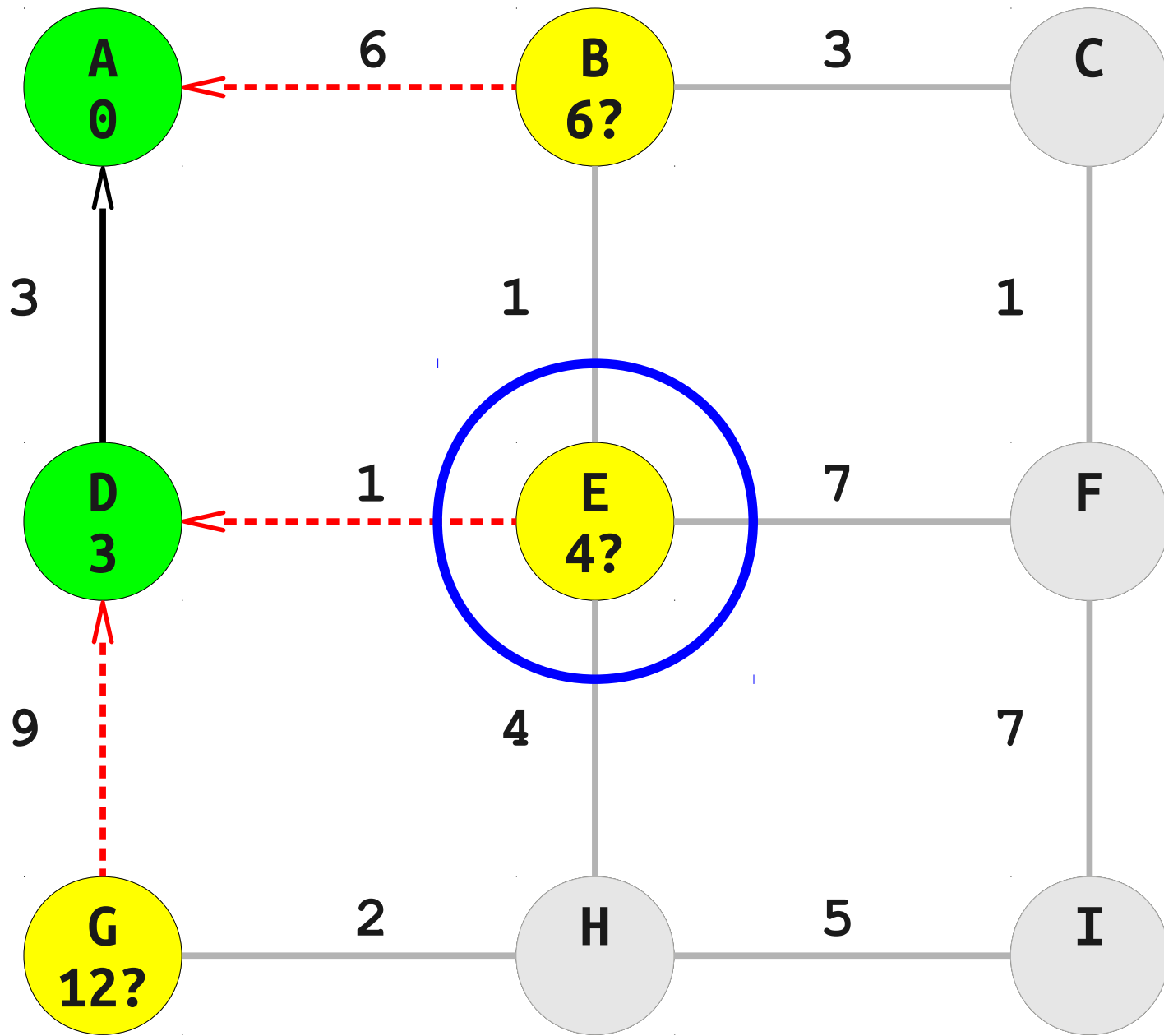
G
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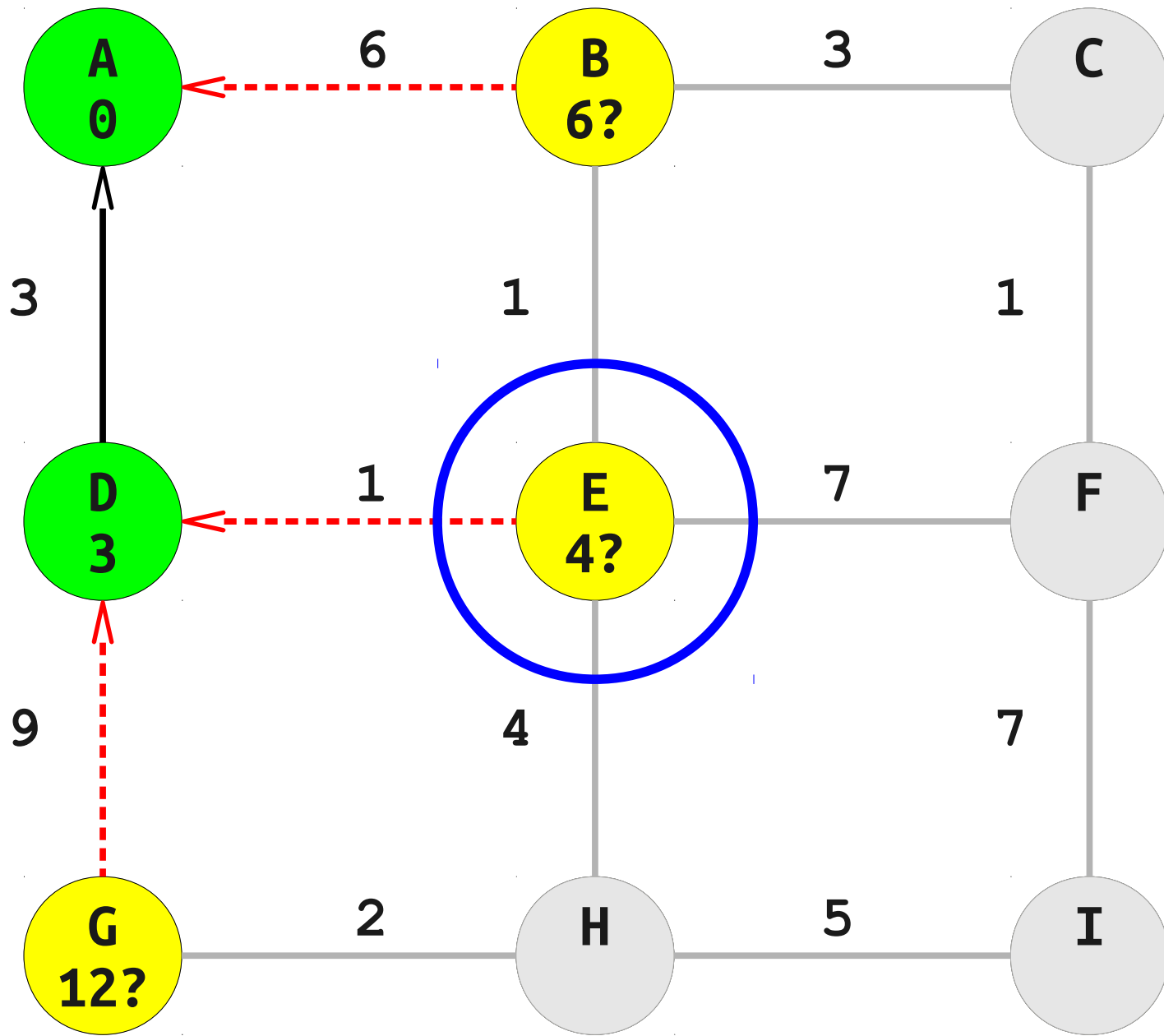
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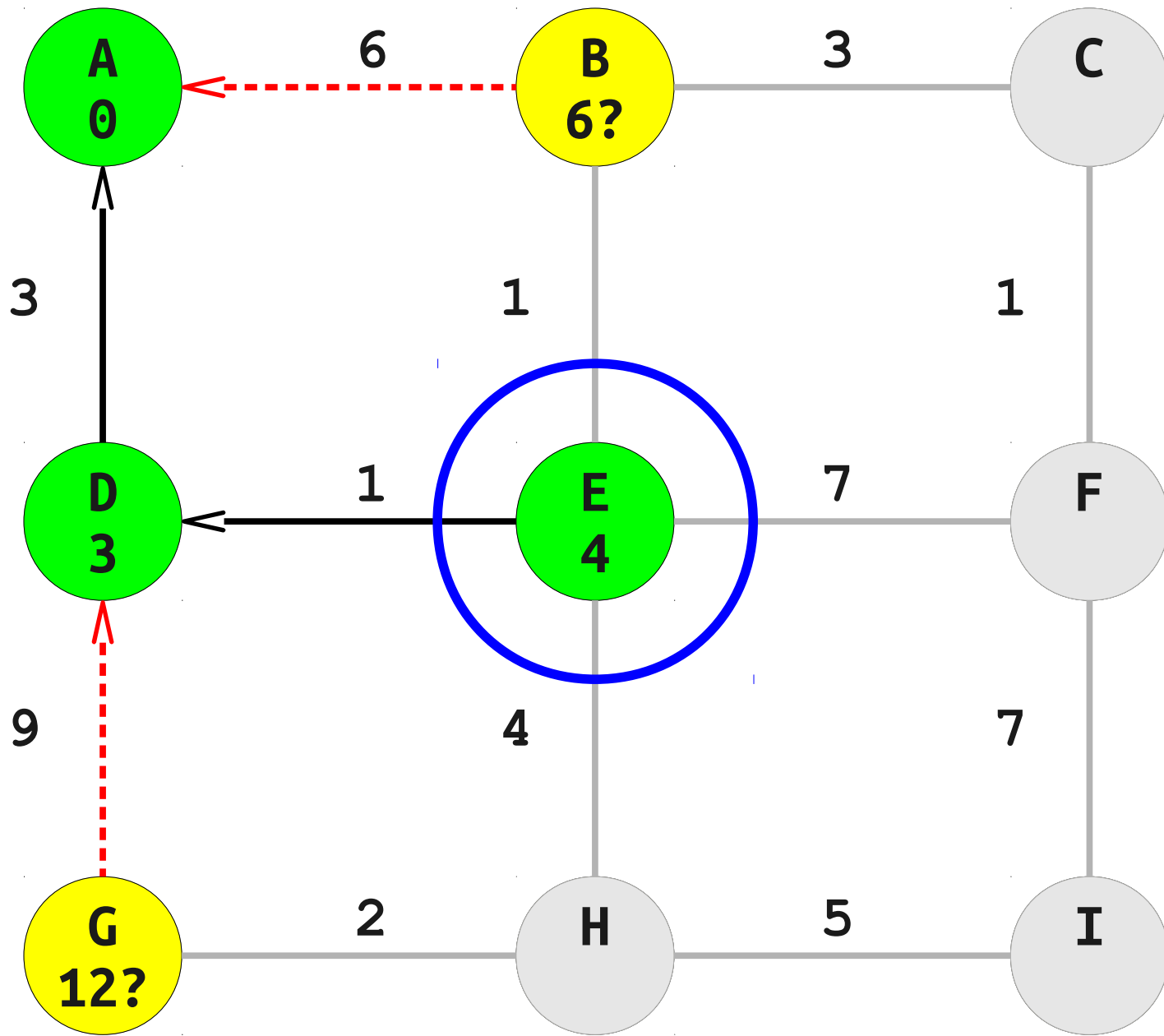
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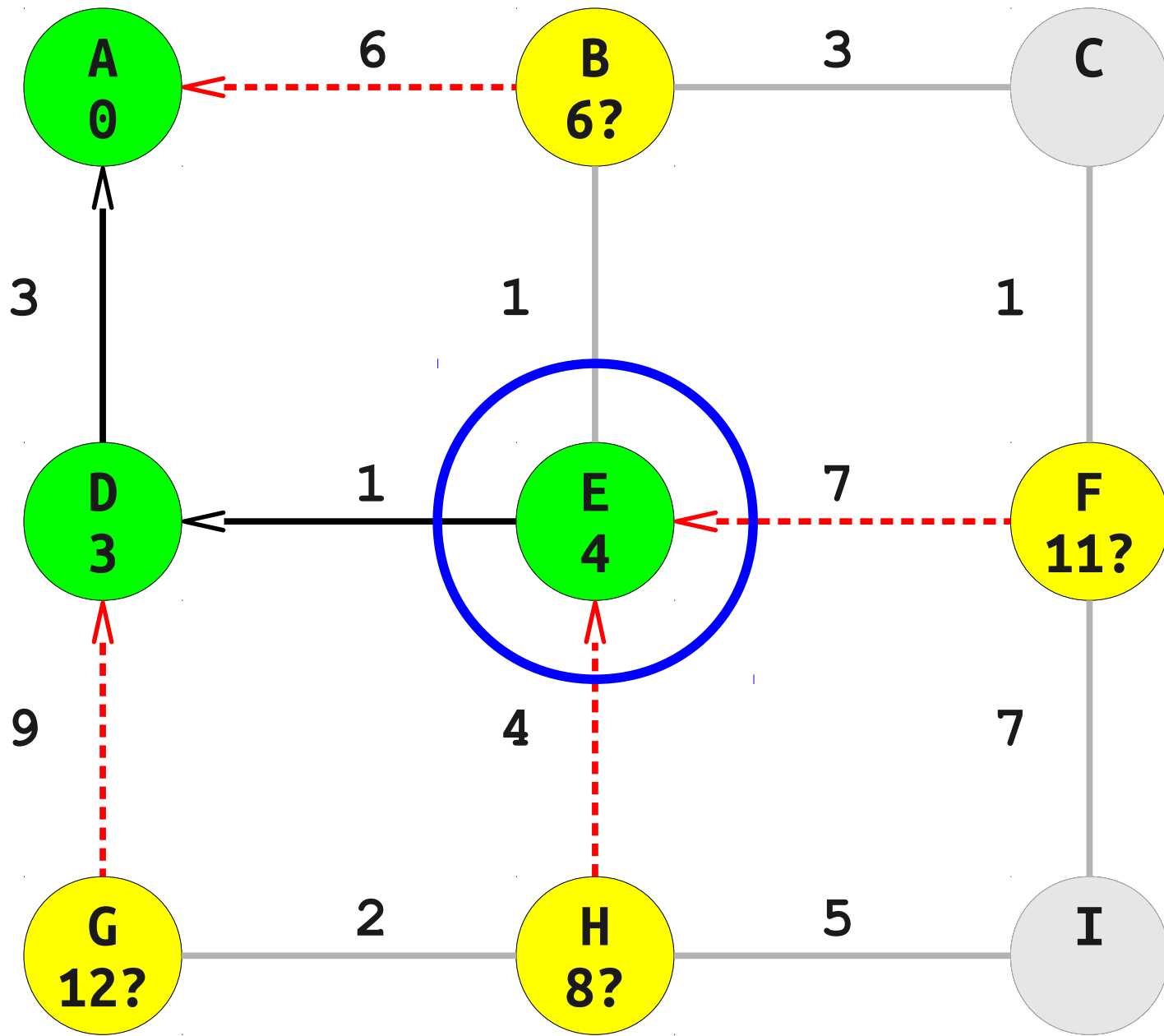
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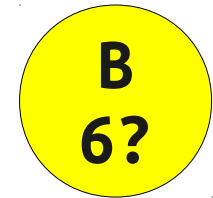
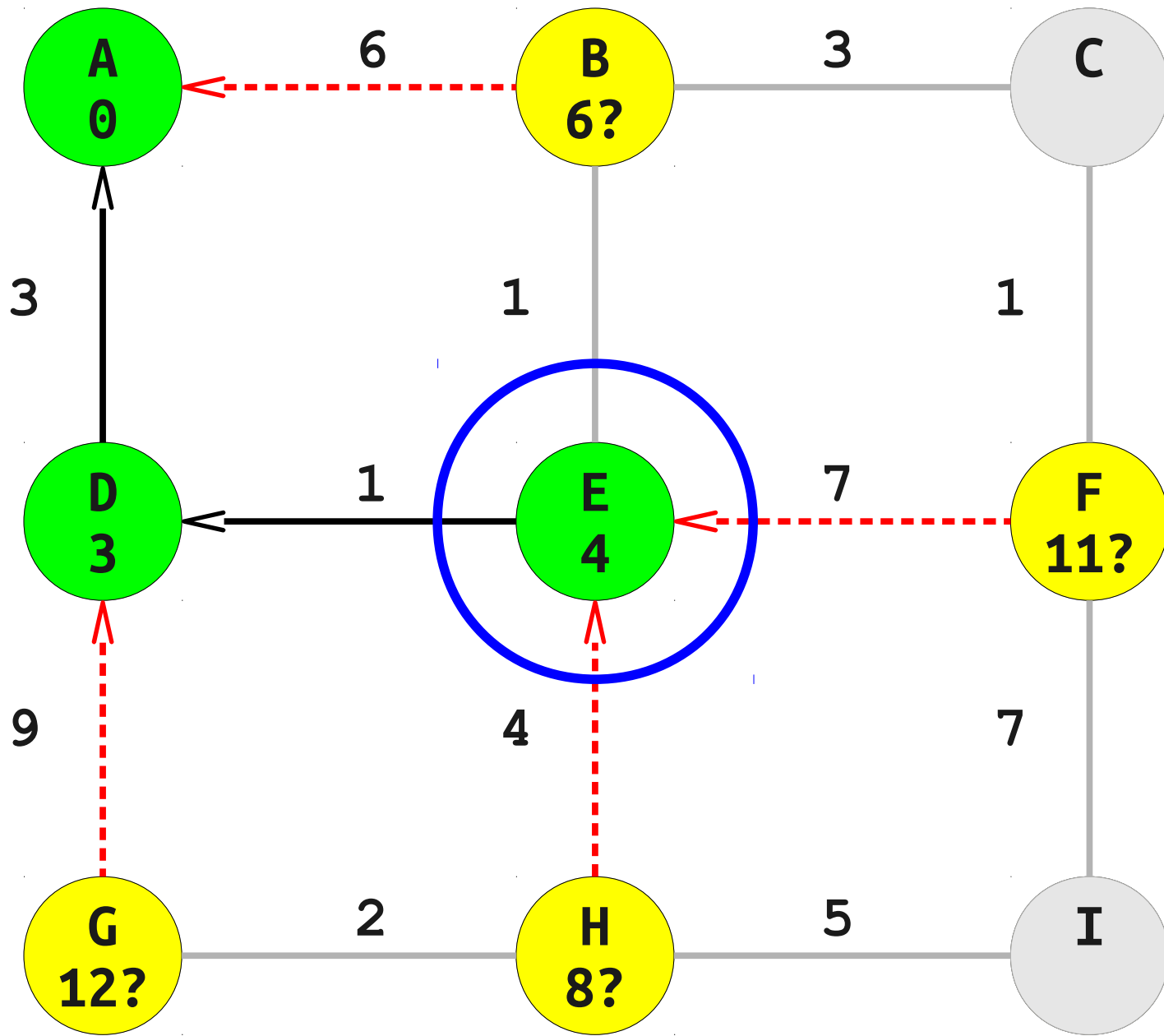
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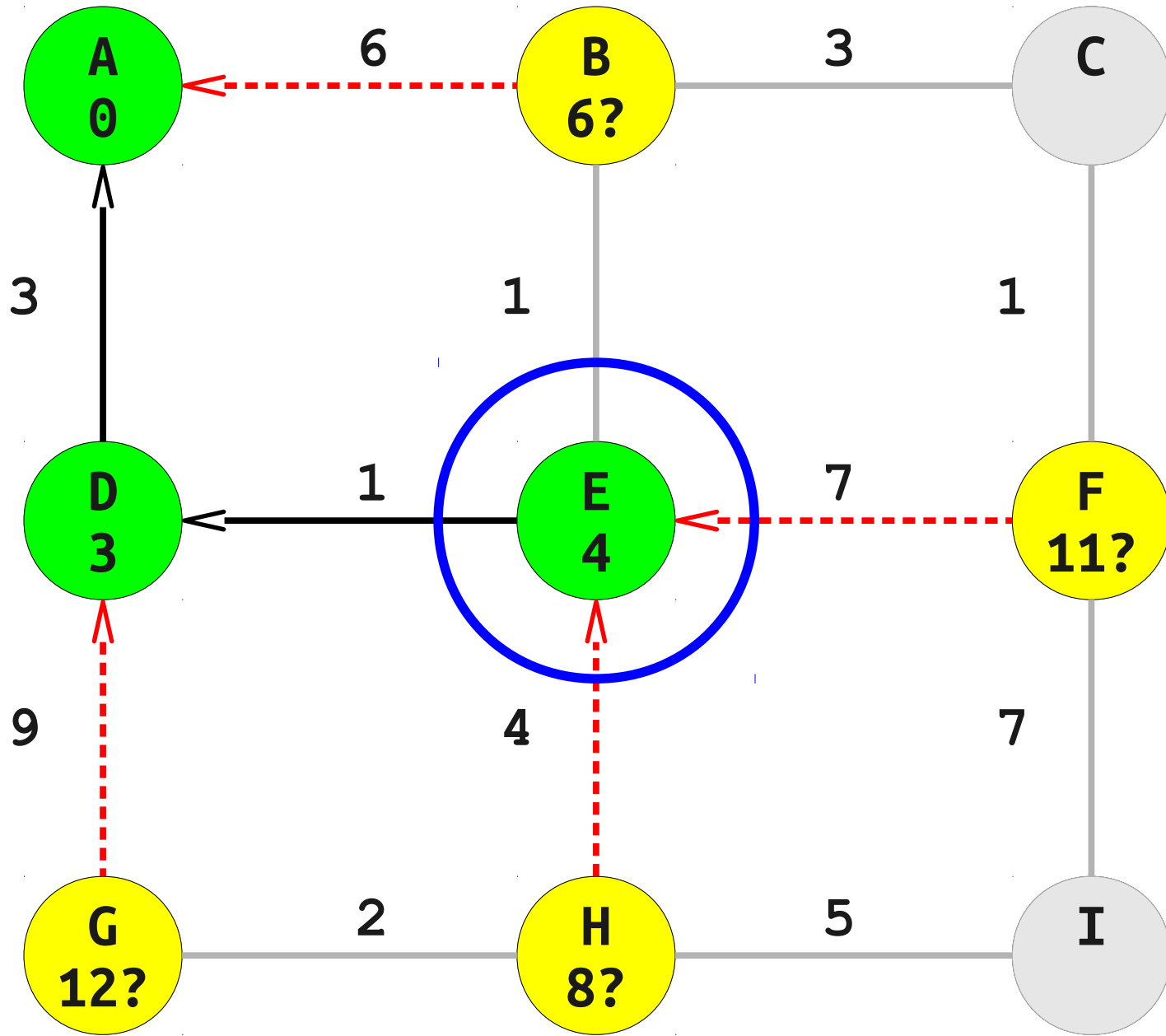
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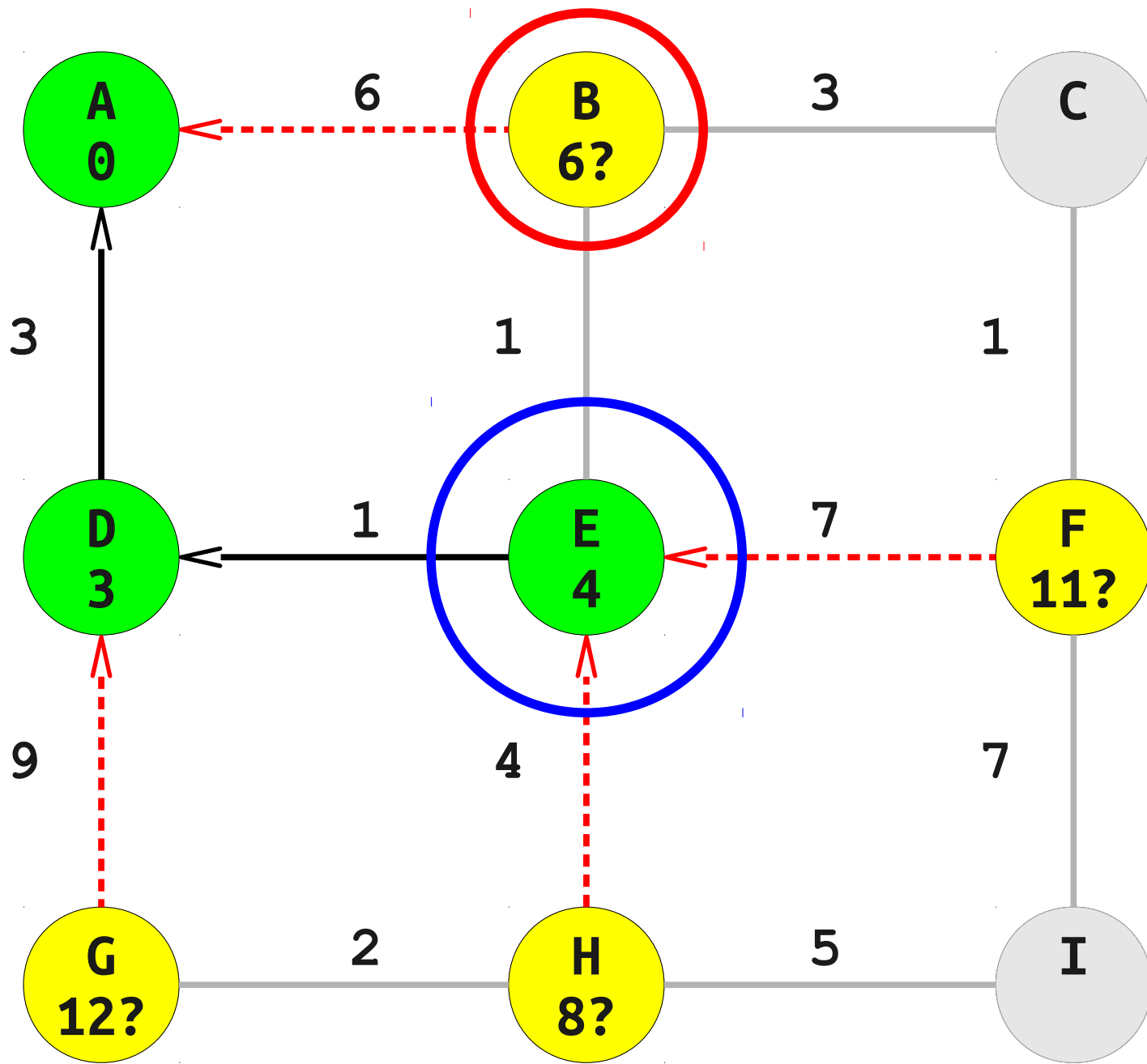
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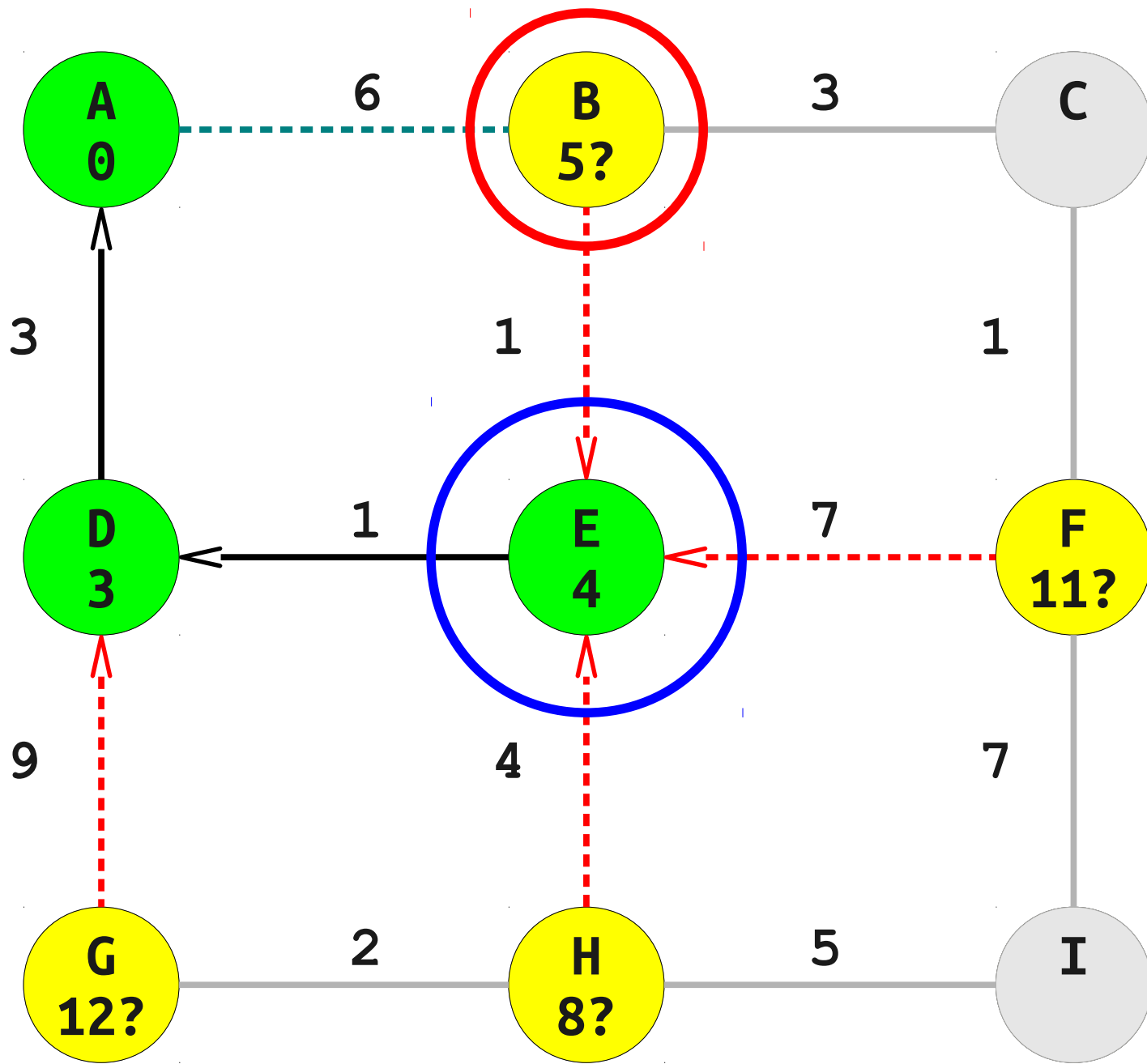


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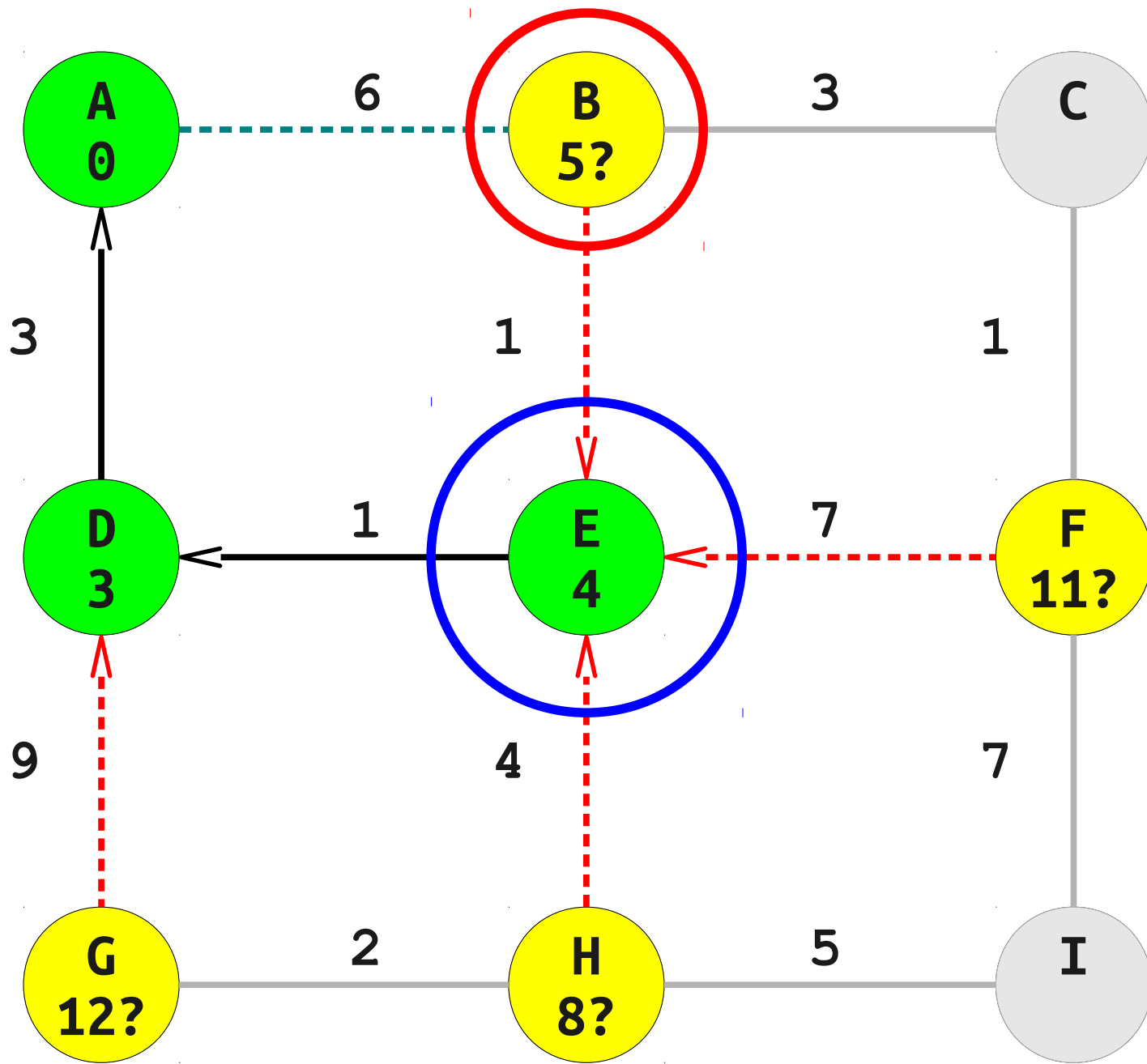
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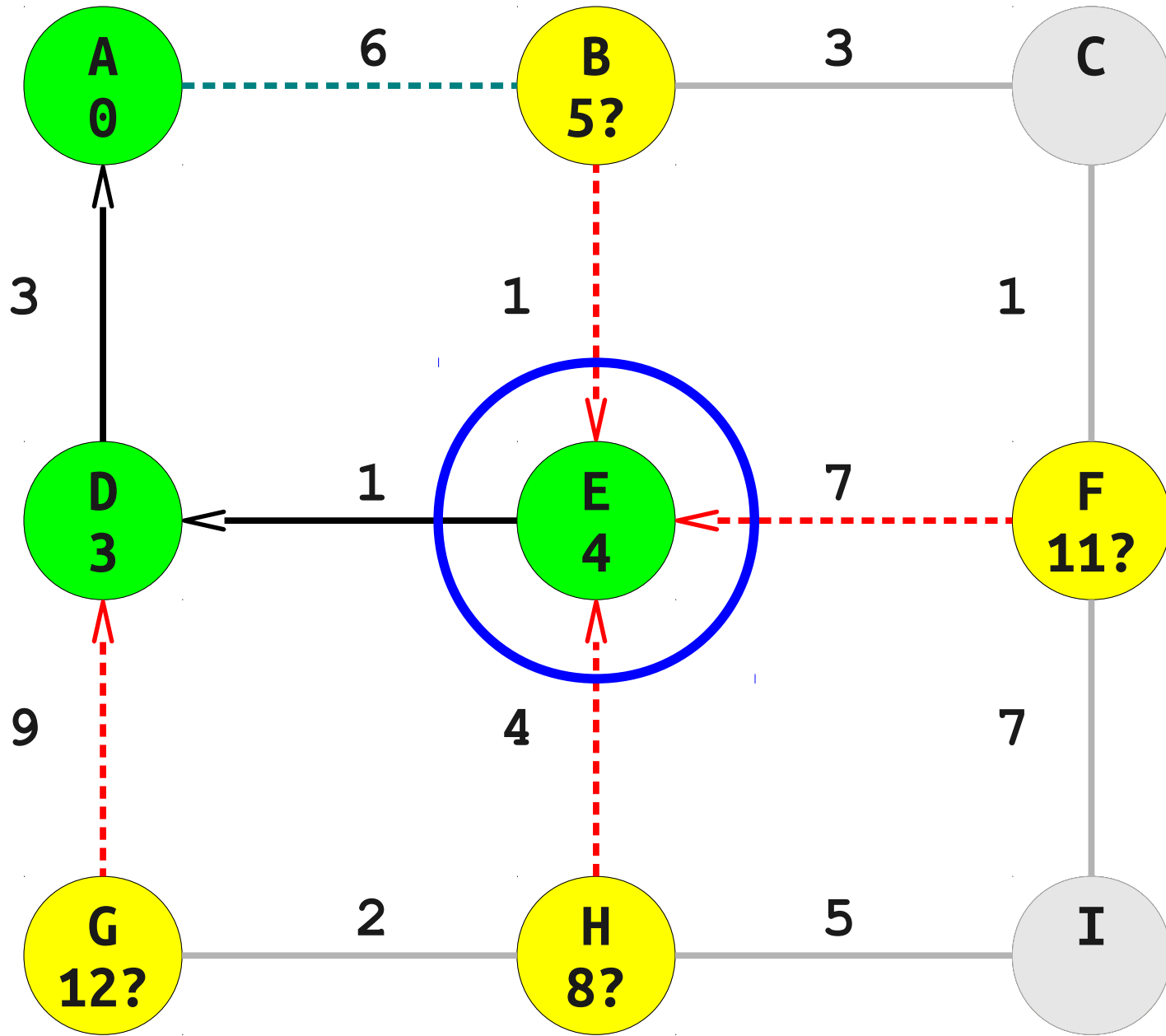
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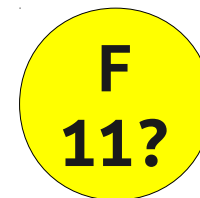
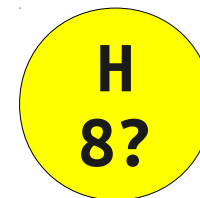
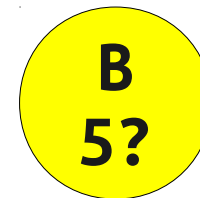
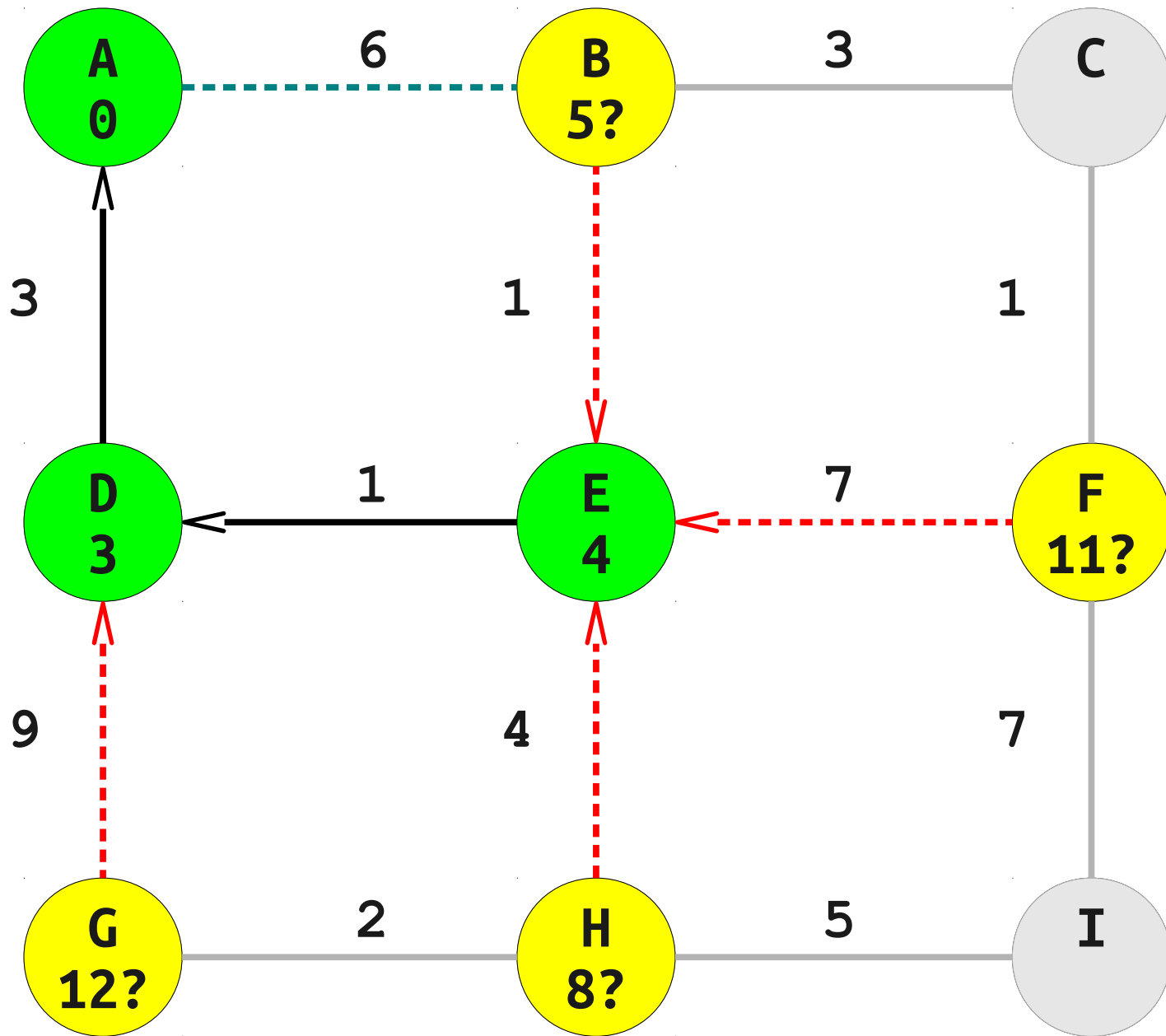


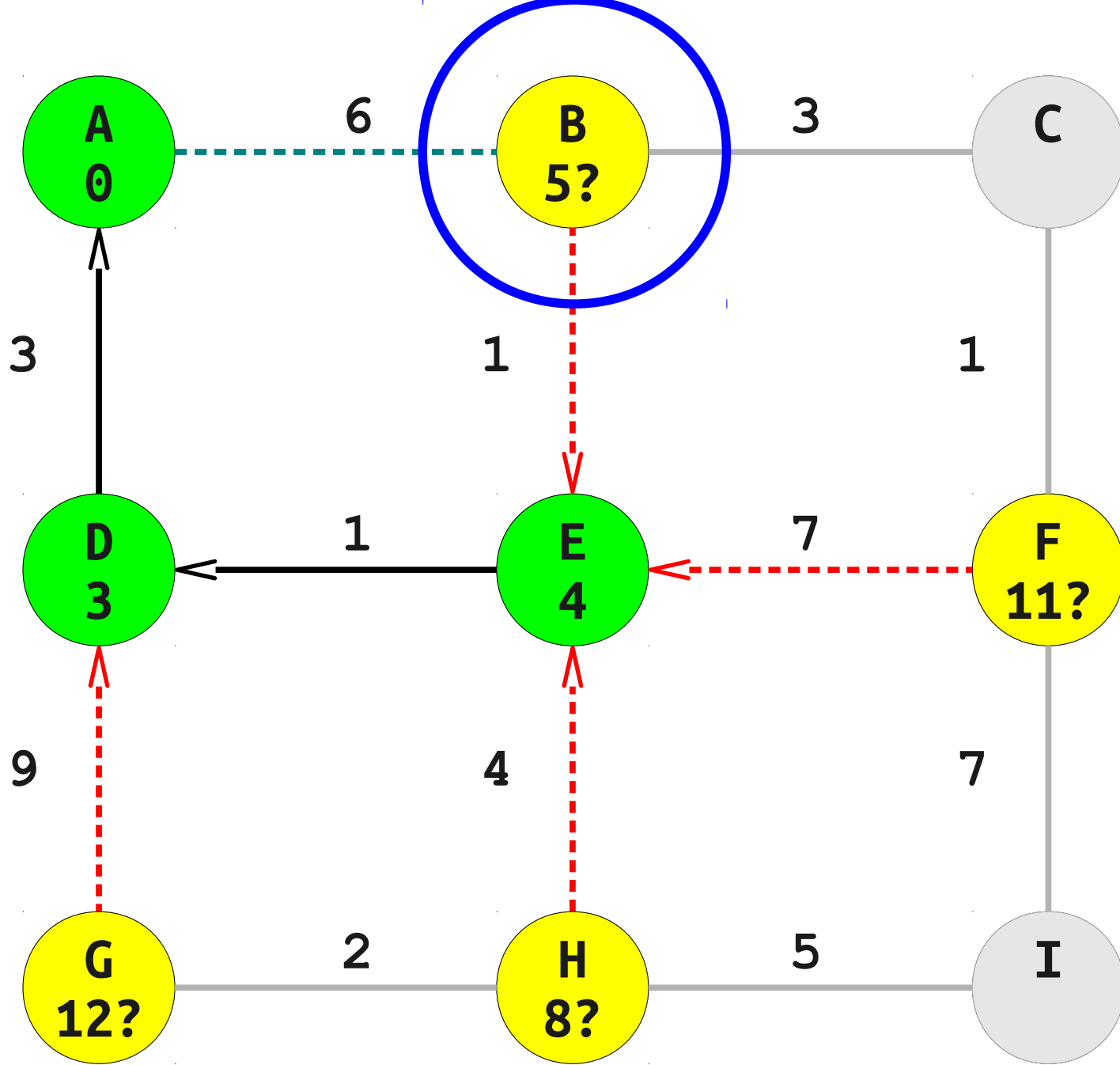
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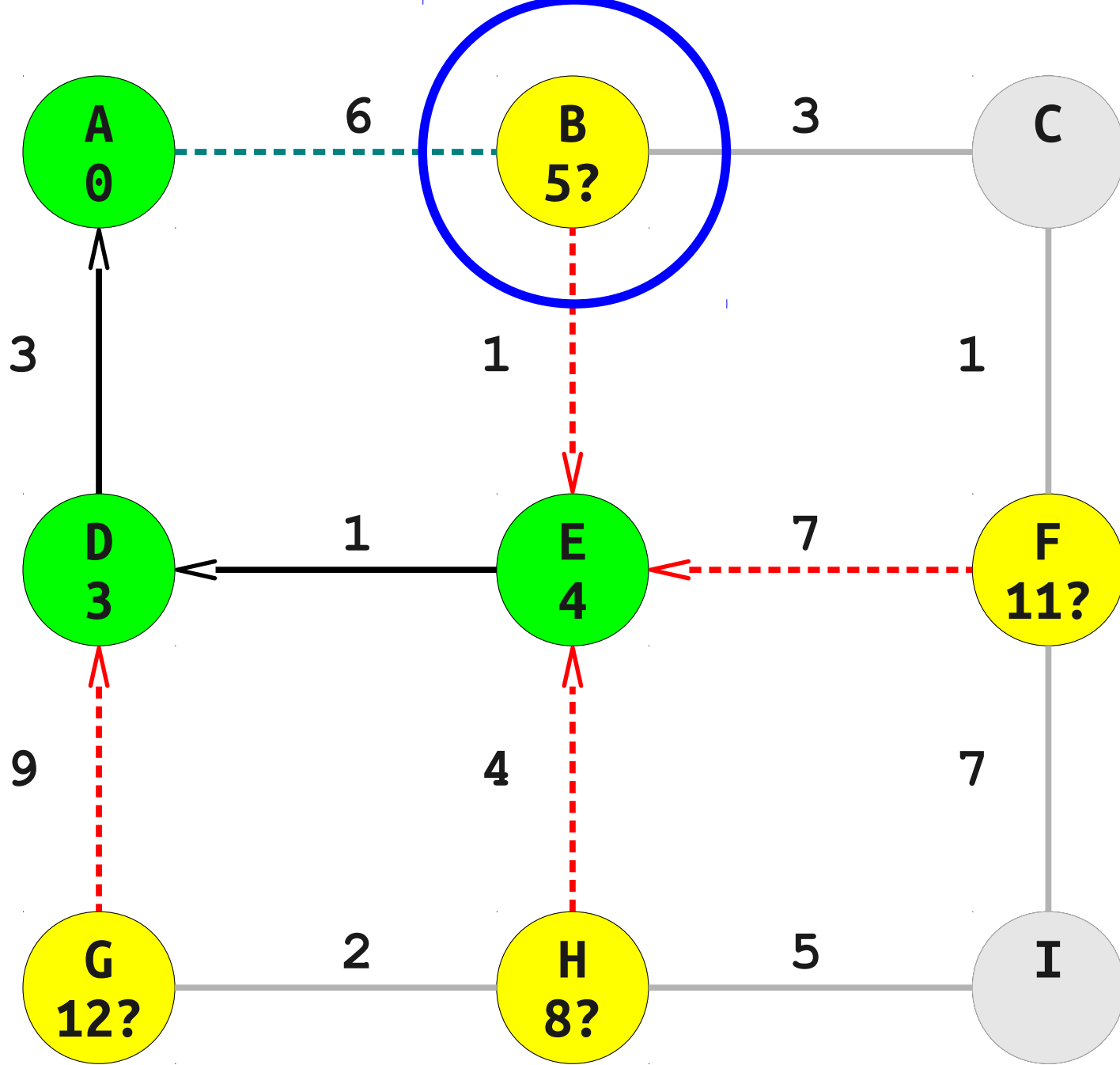


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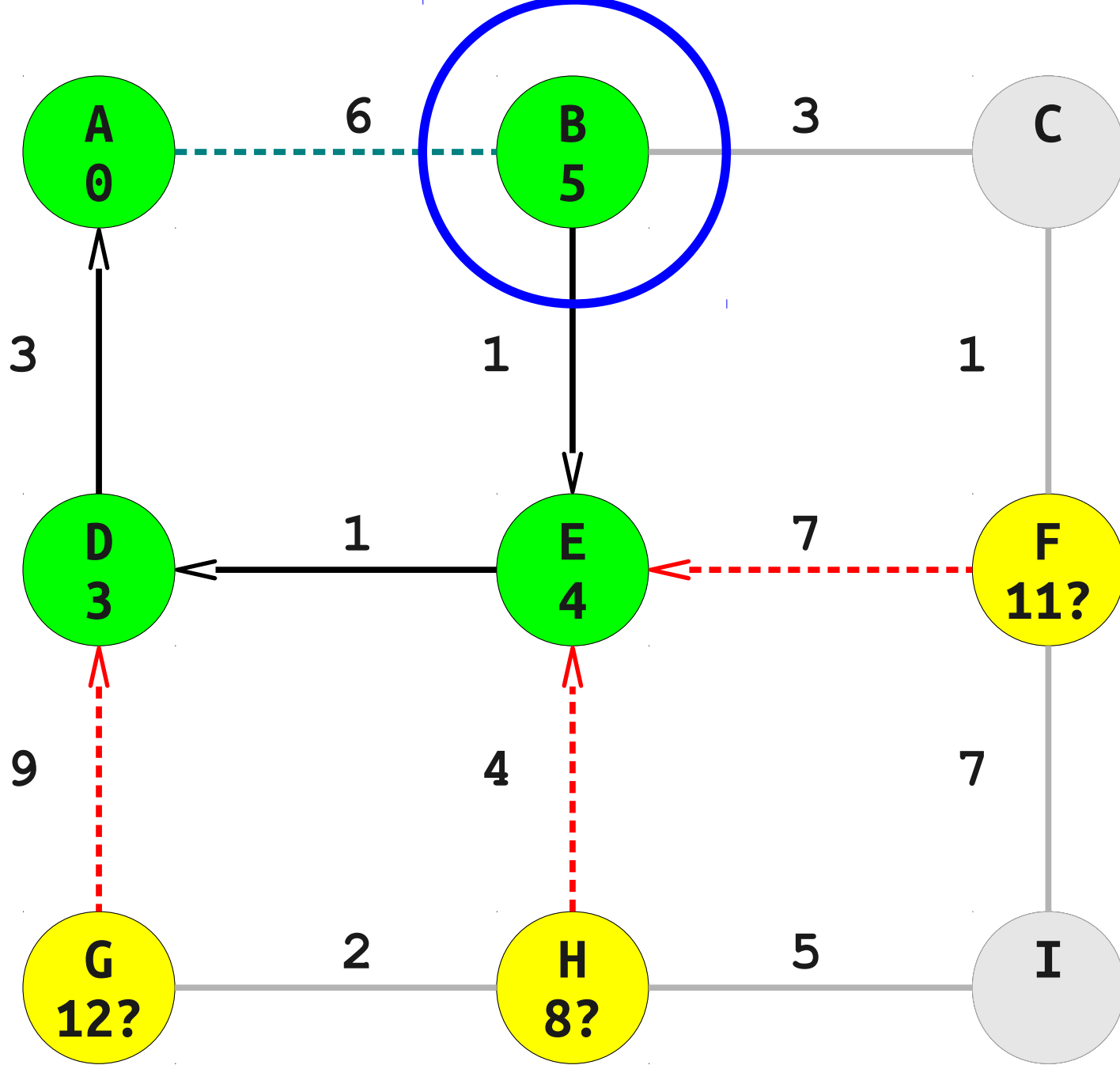
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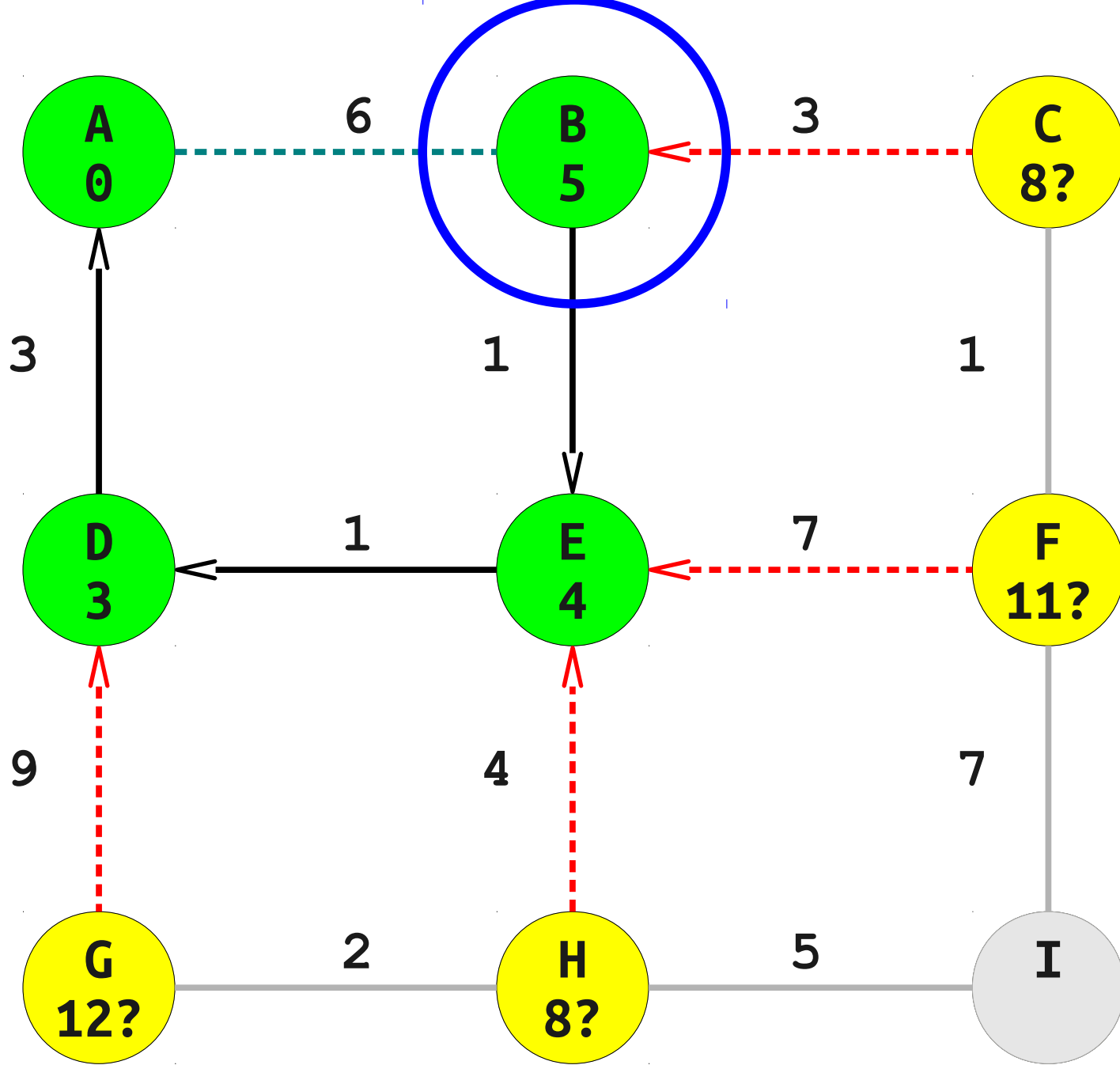
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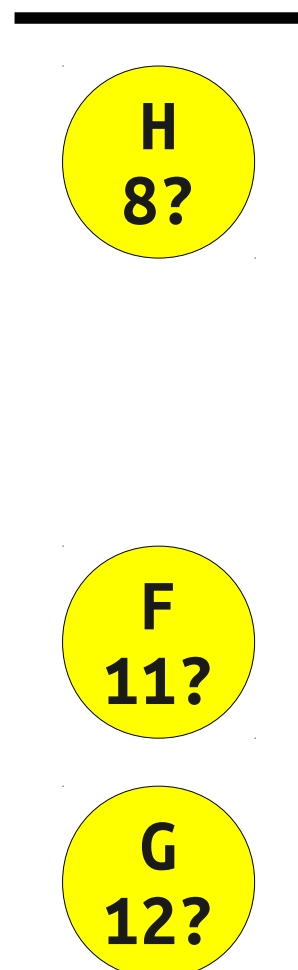
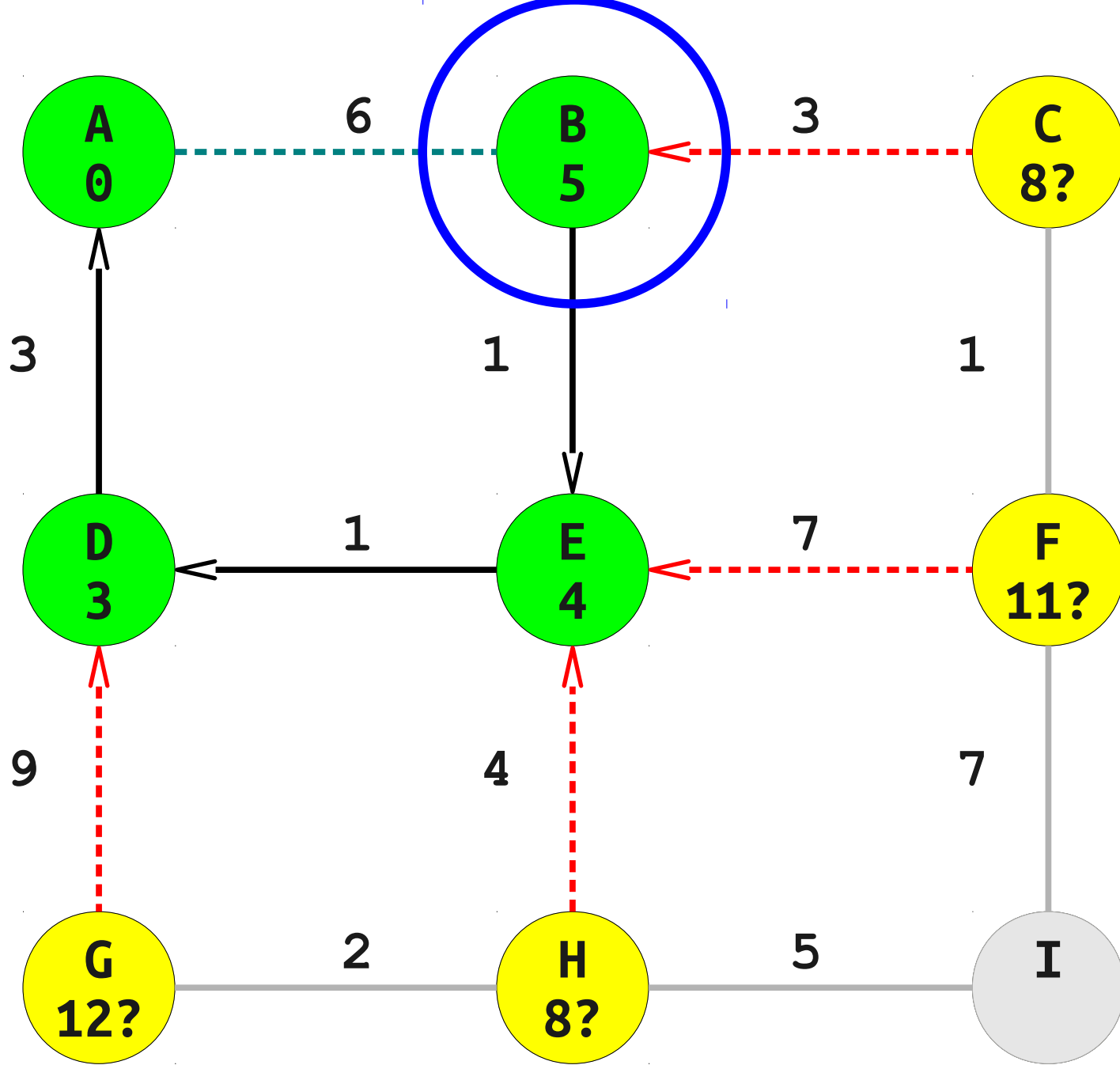
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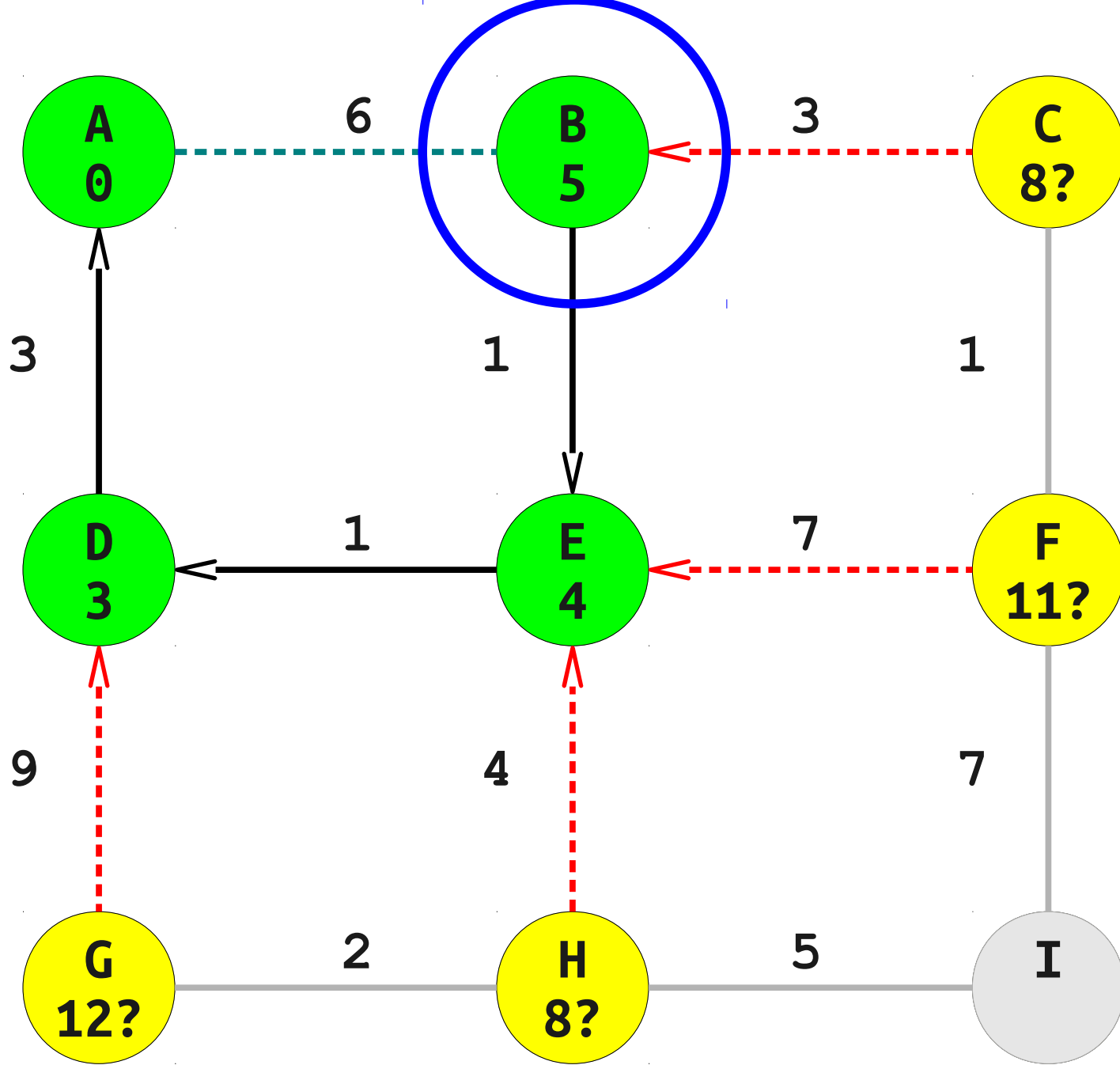


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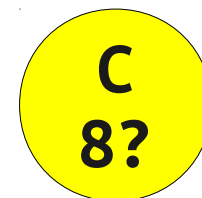
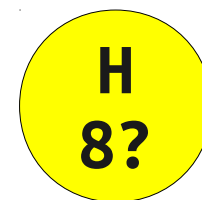
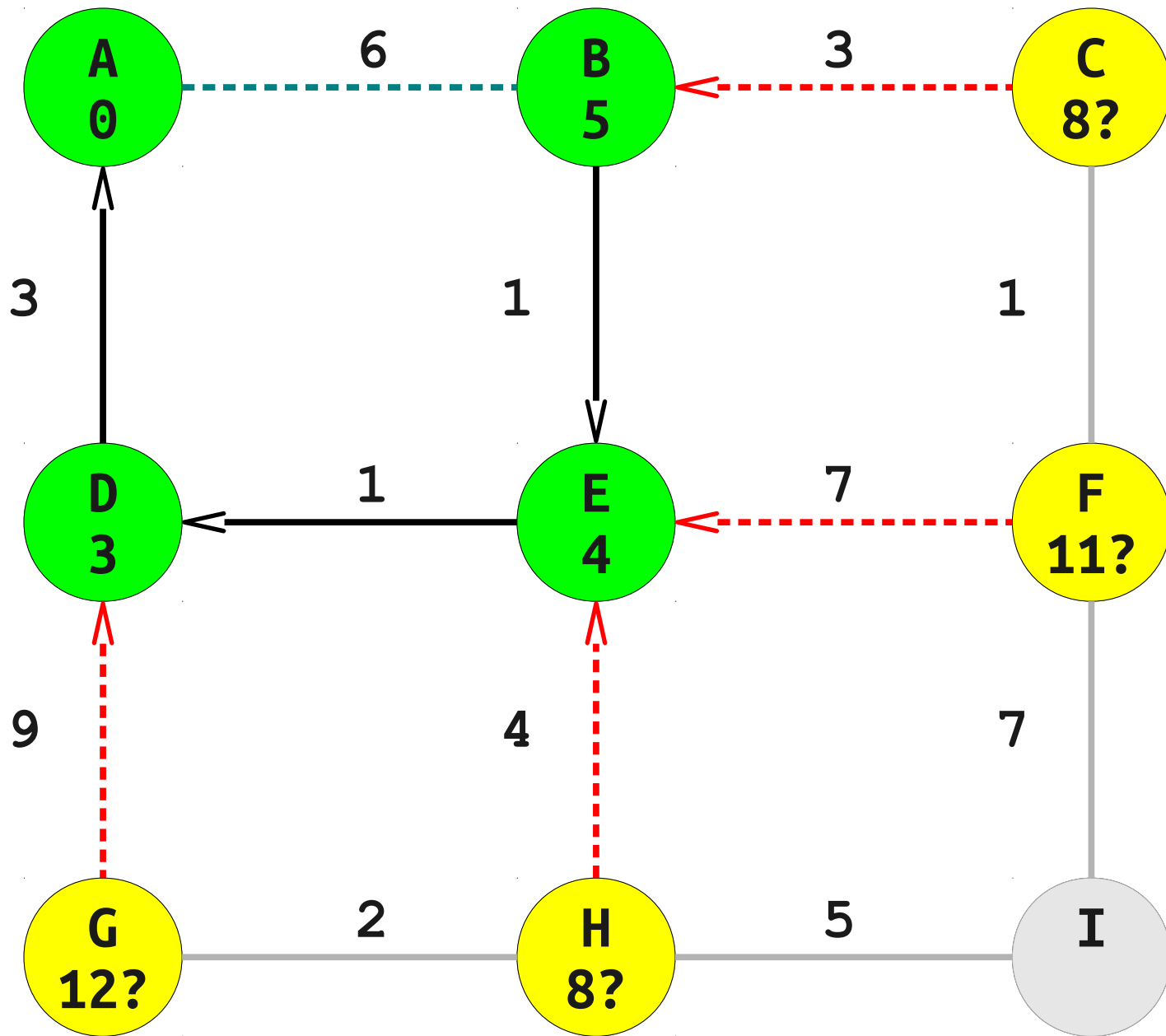
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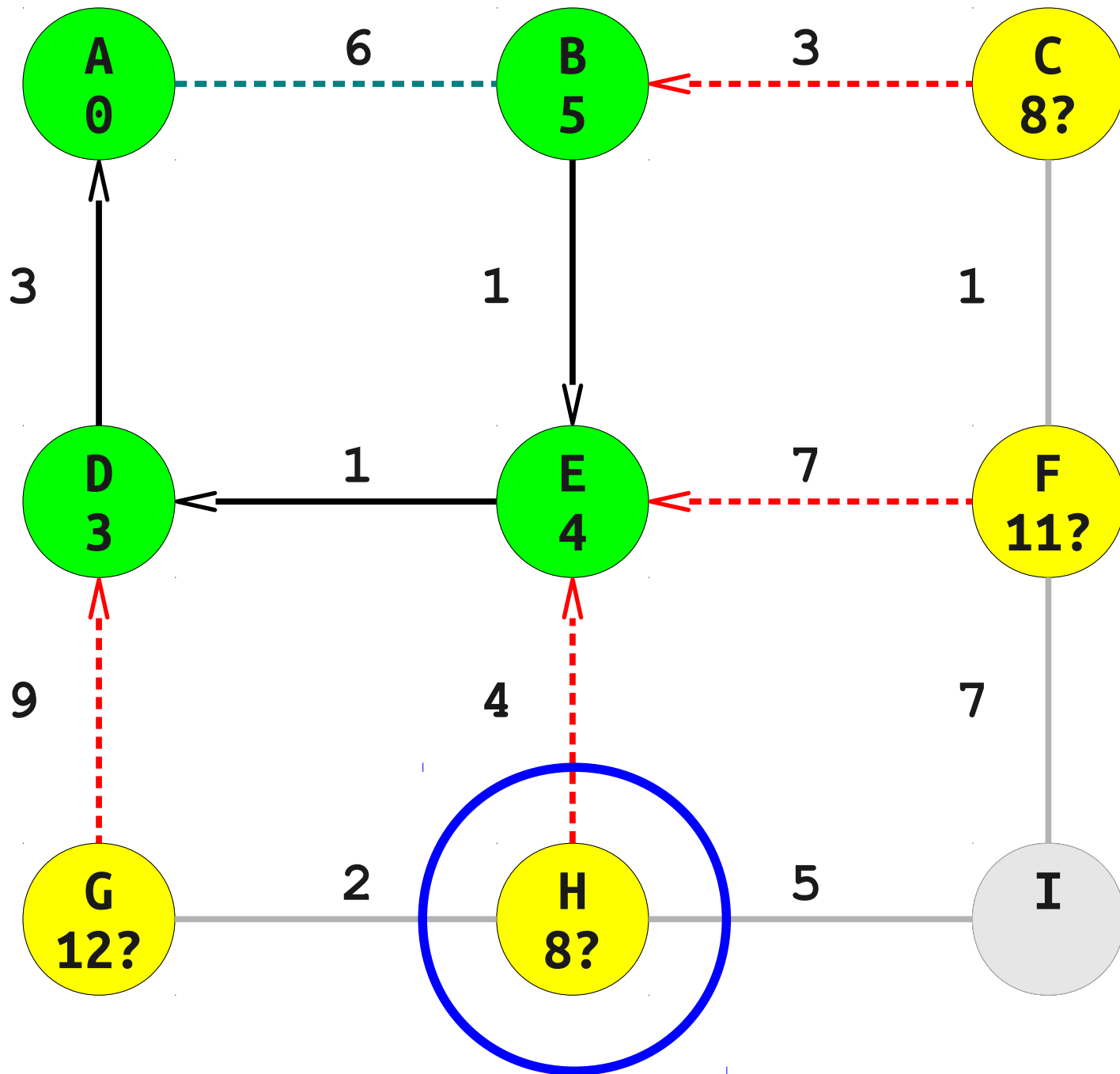
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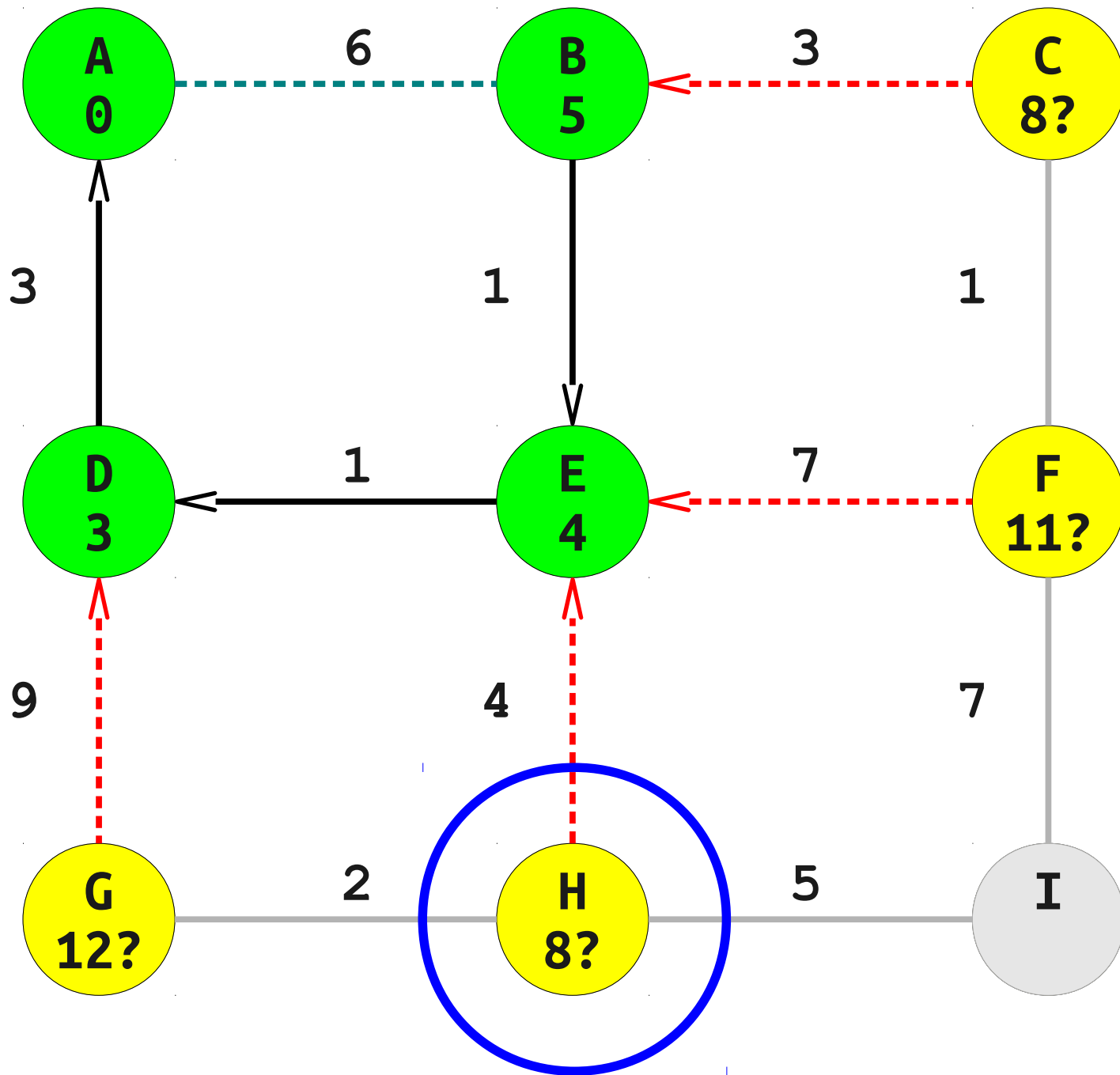


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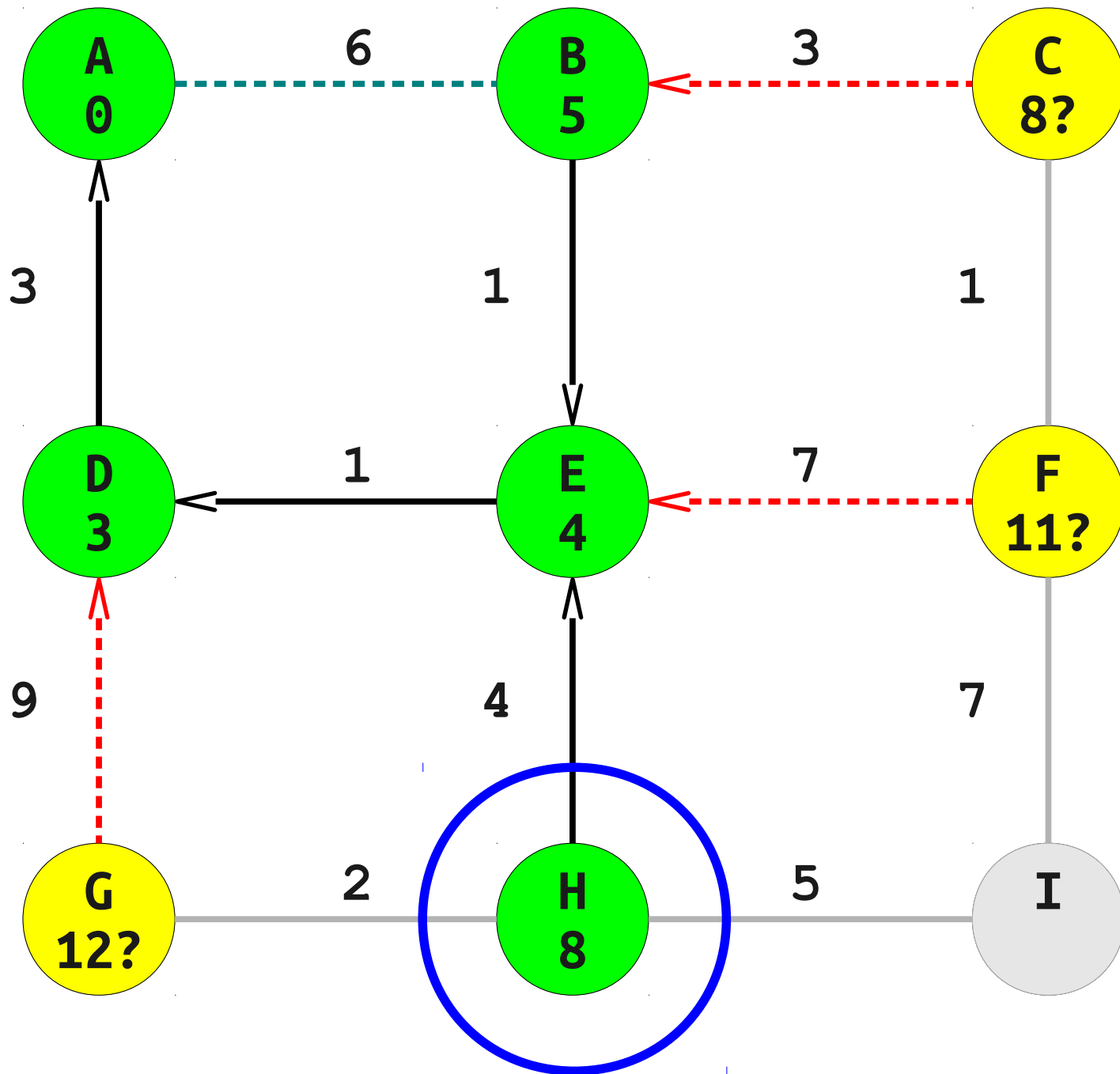
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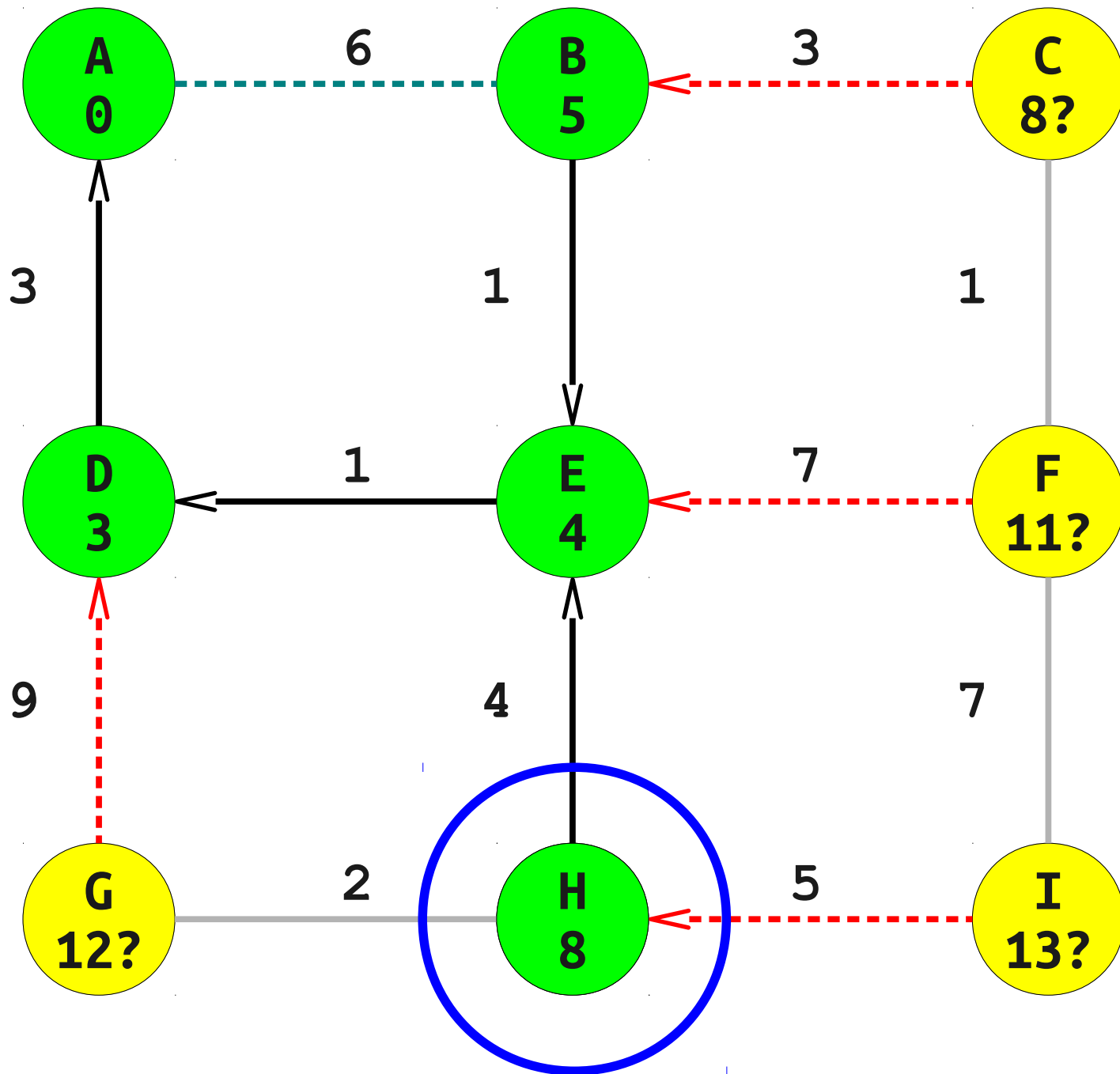
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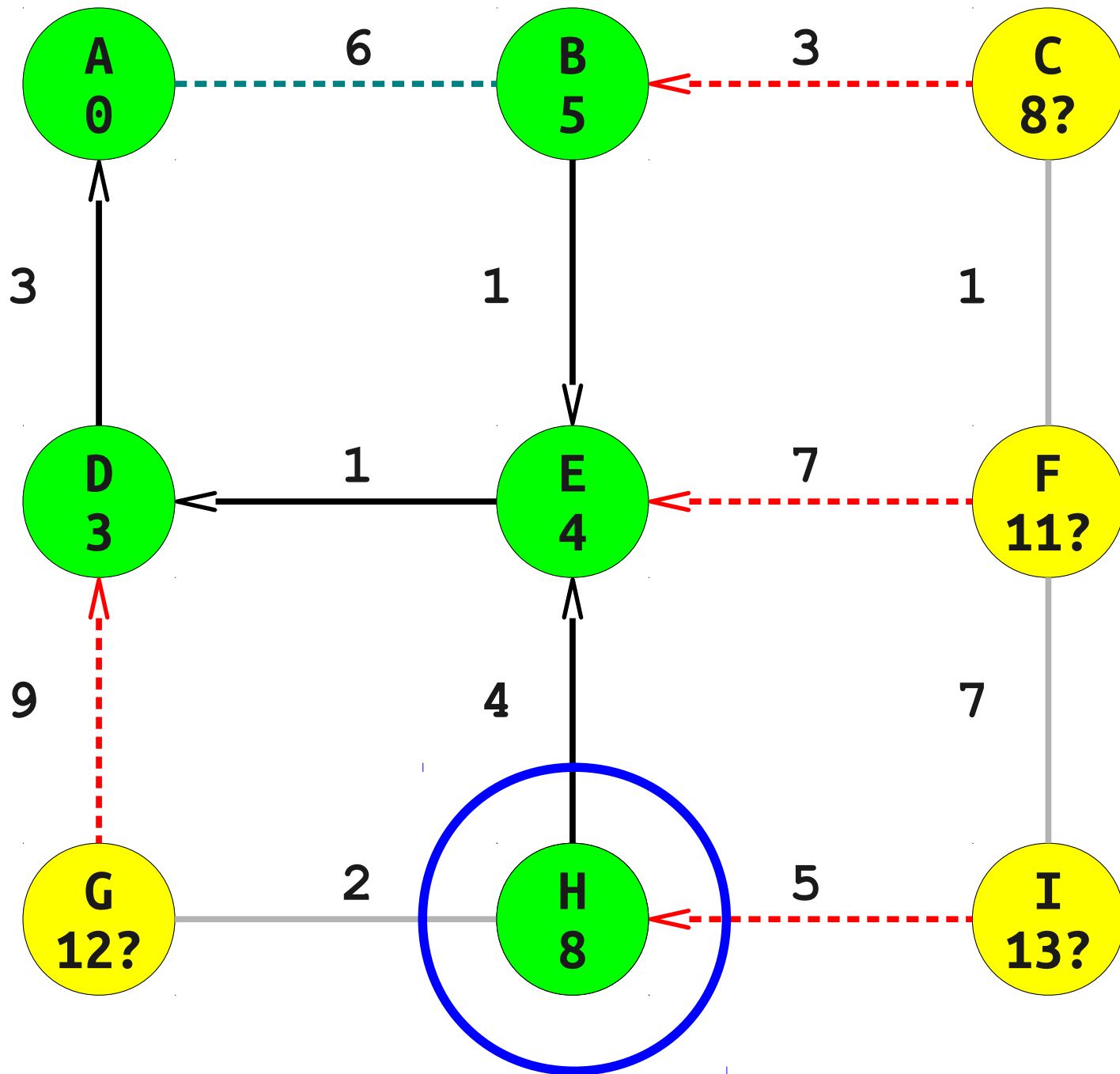
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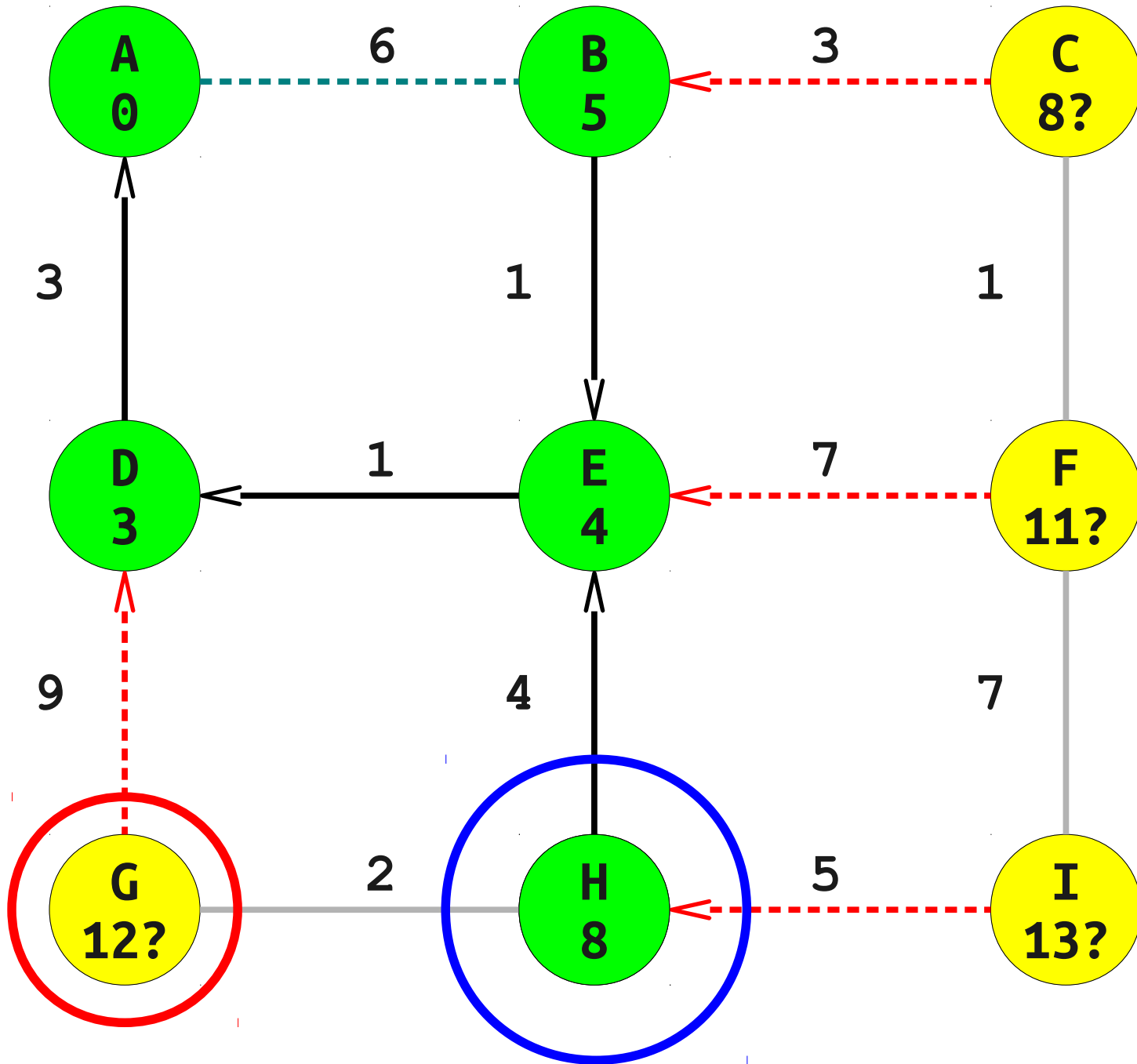


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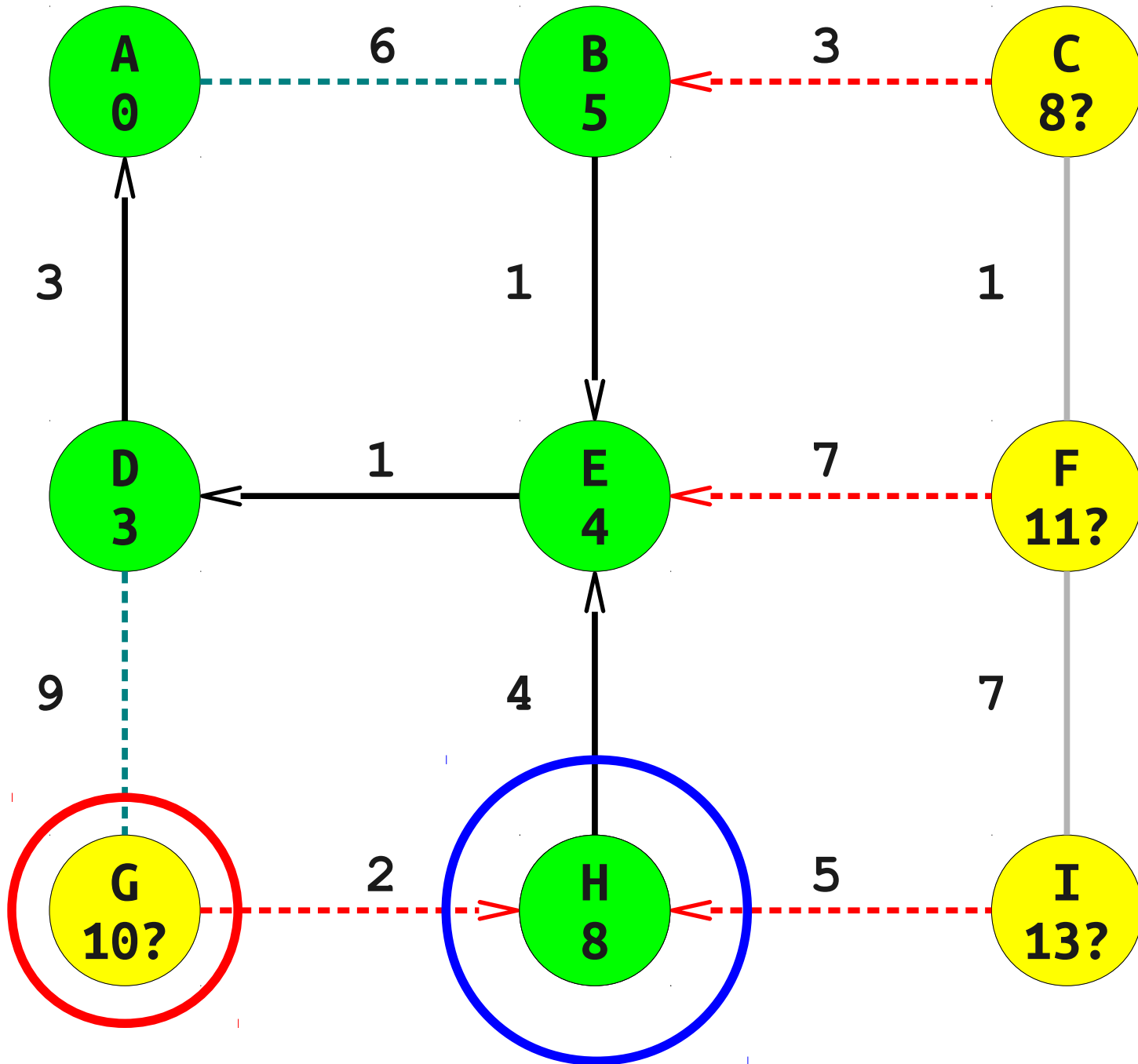


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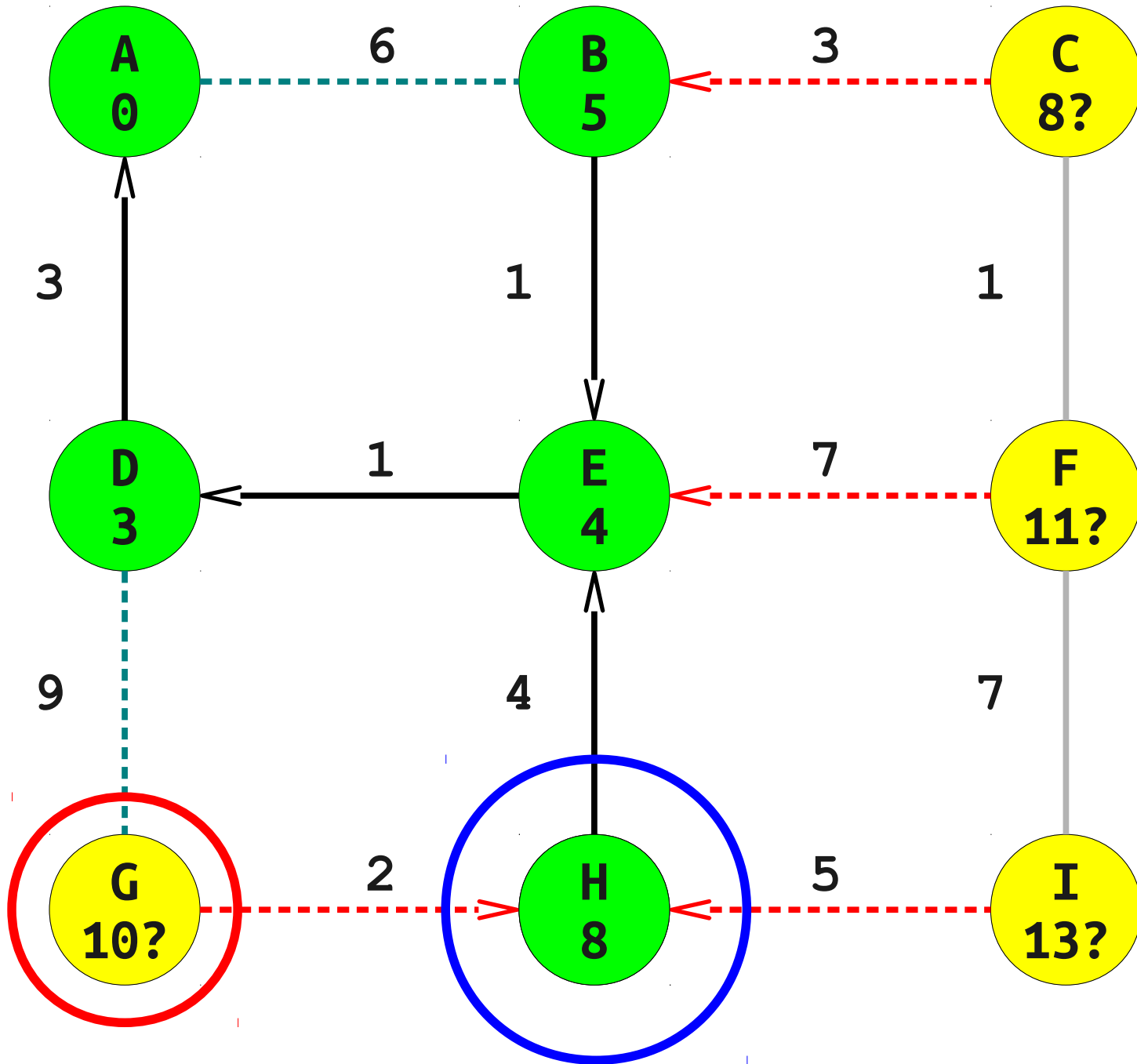
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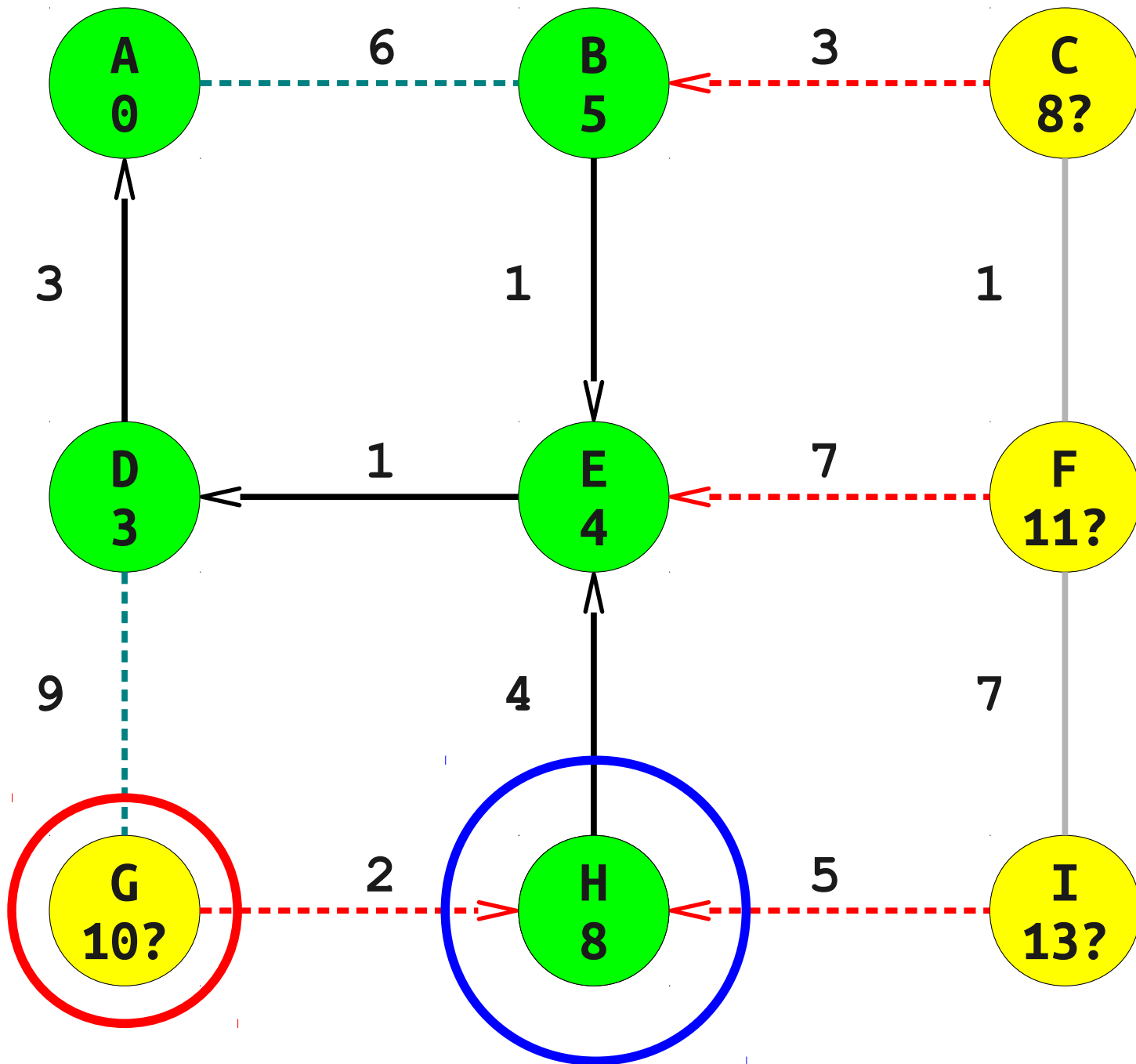


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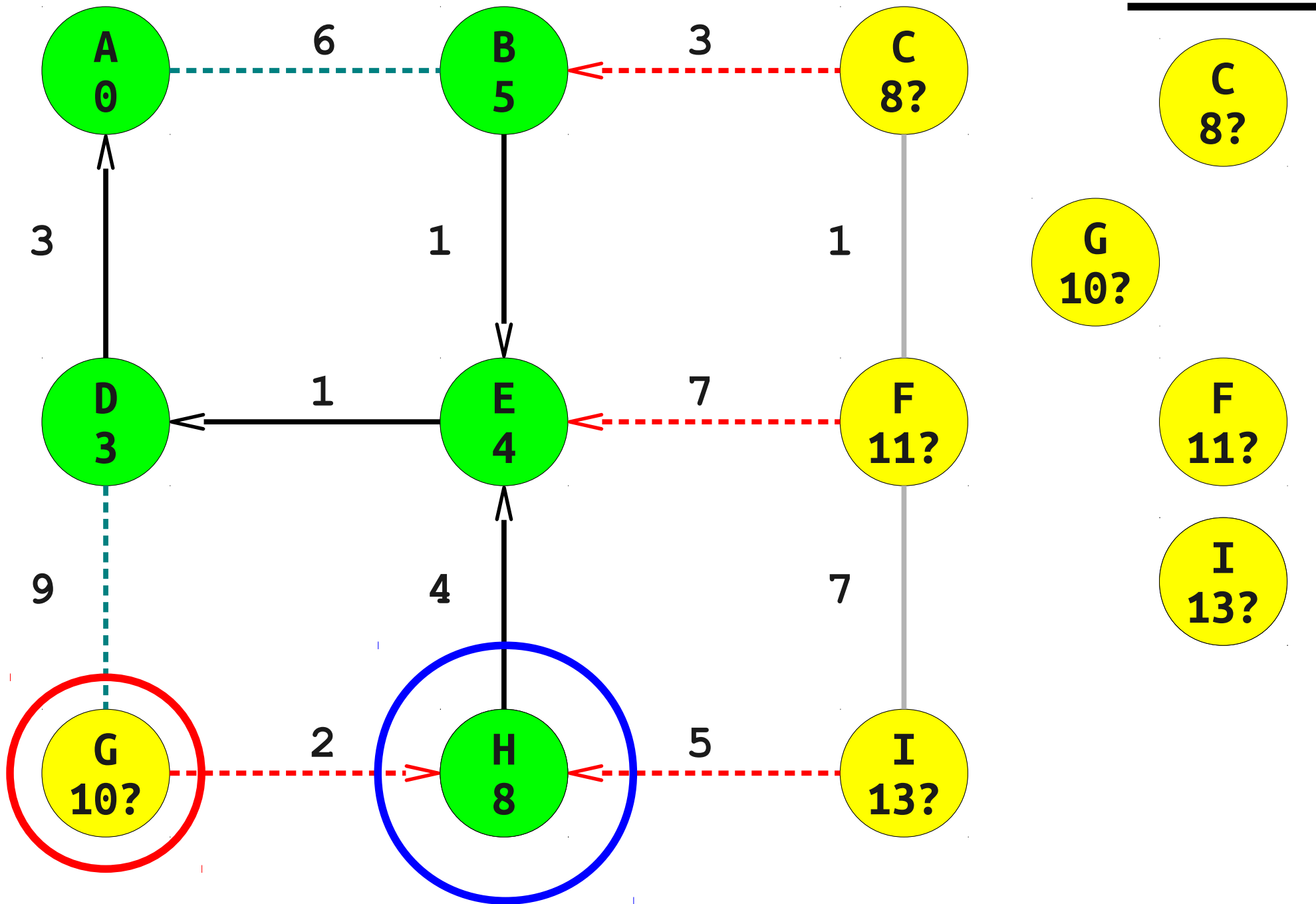


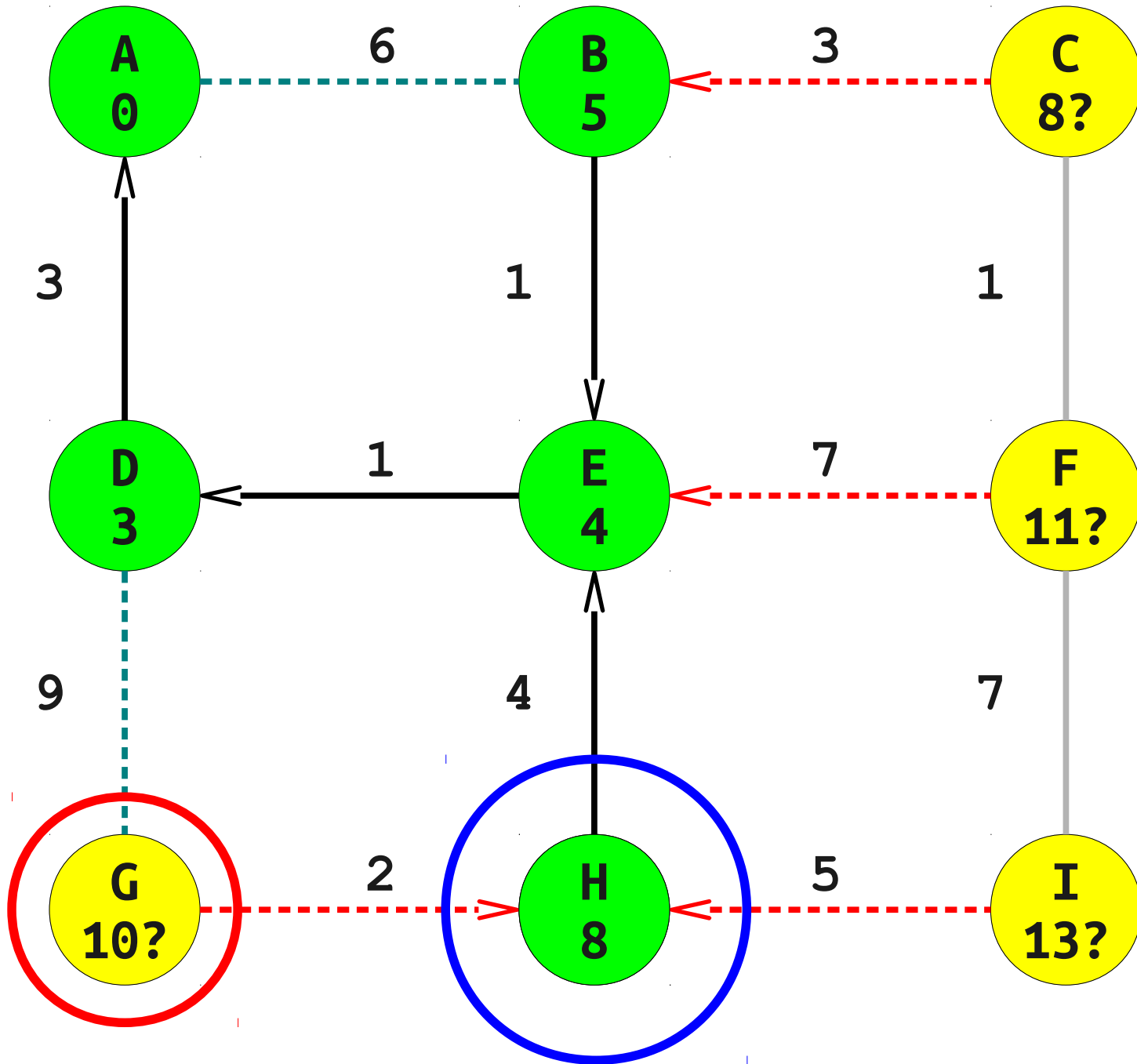
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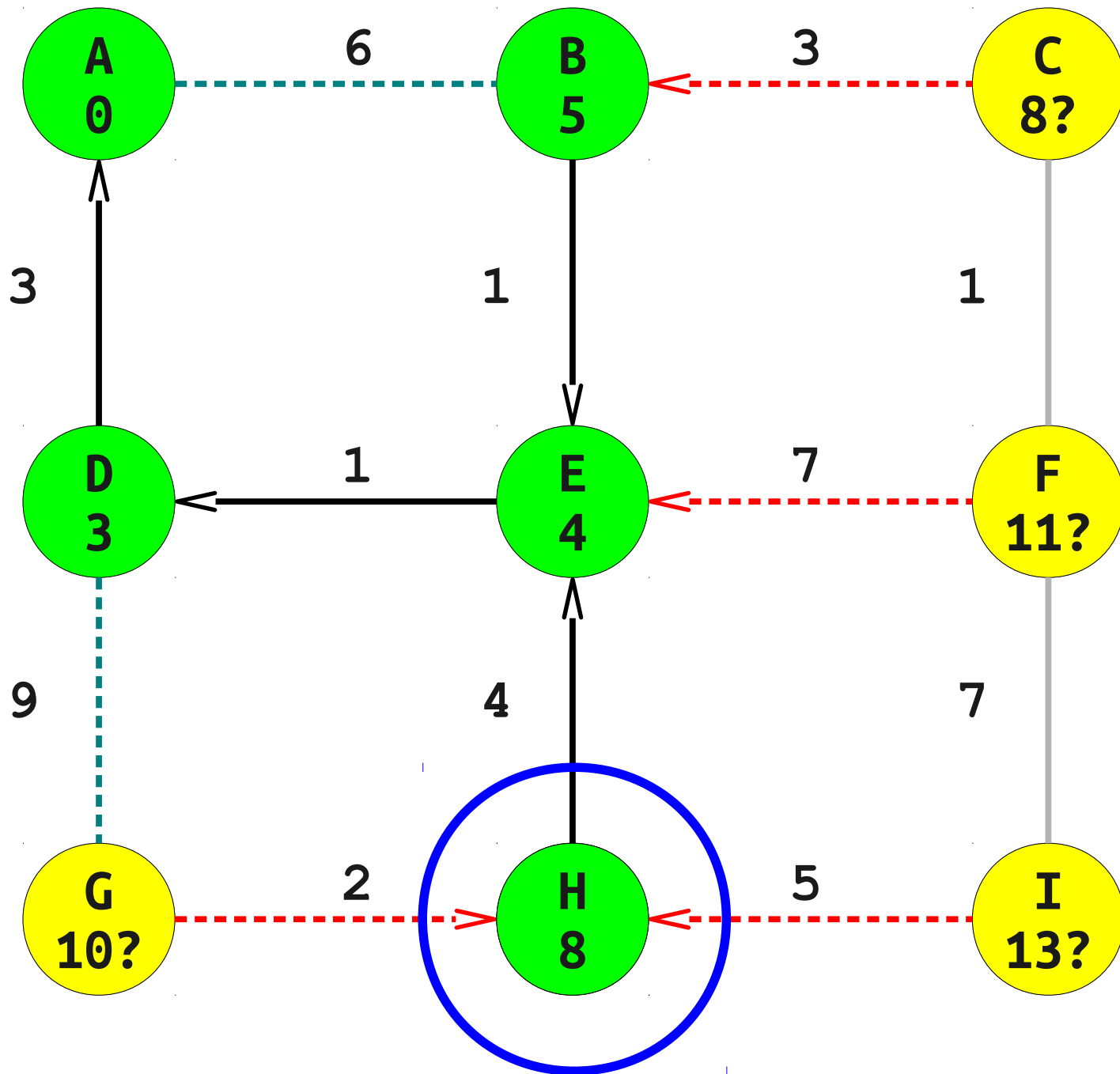


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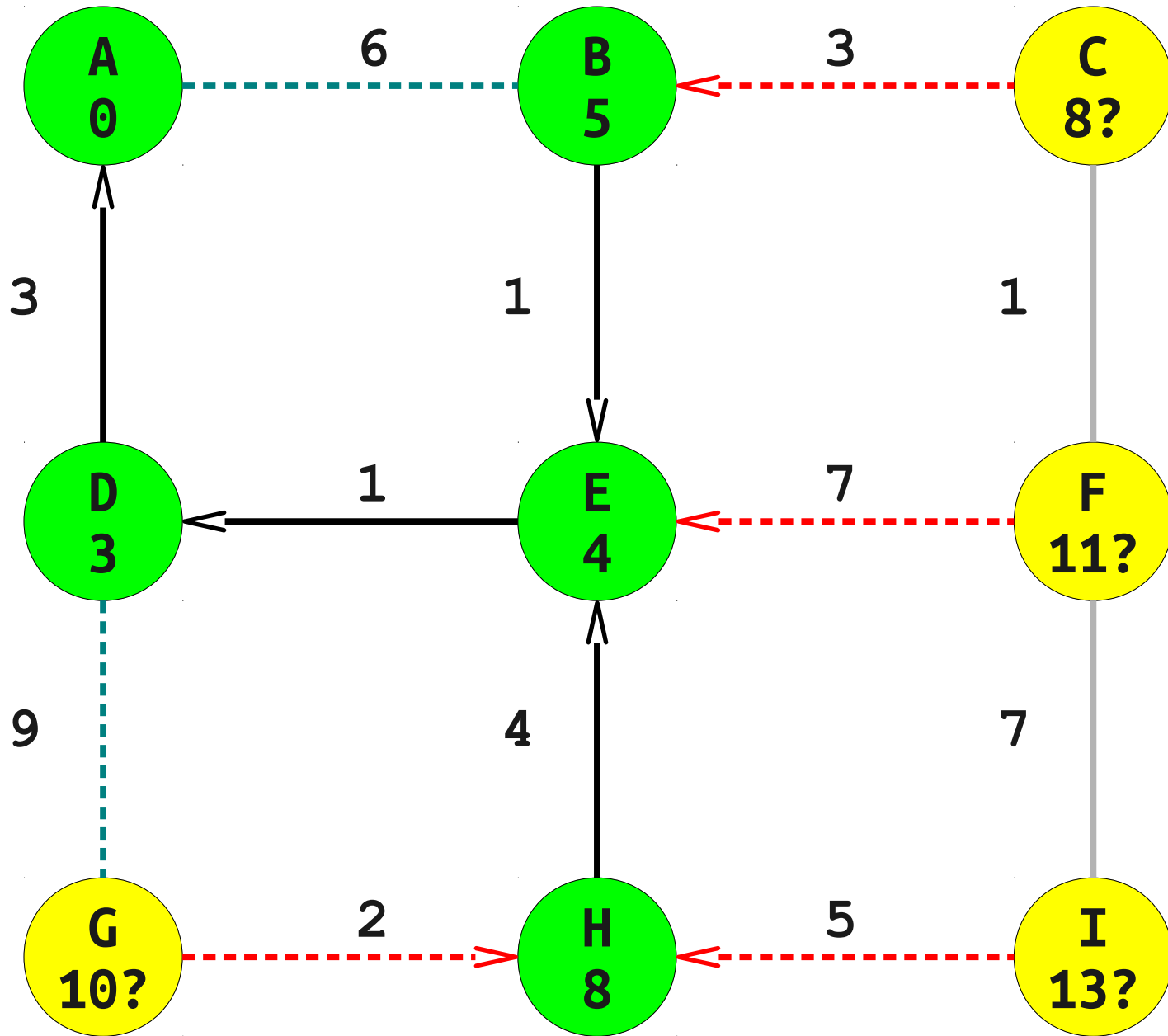
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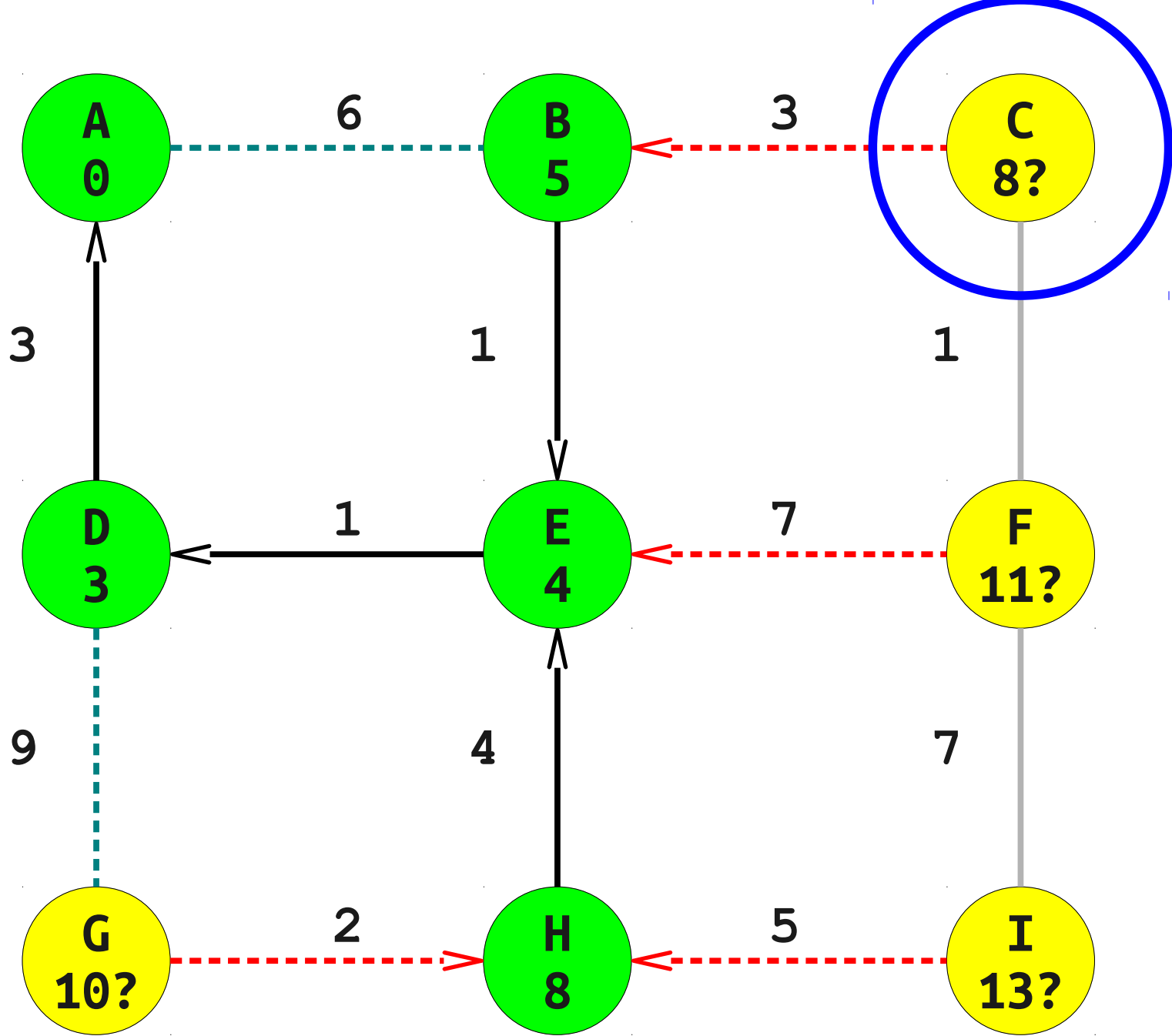


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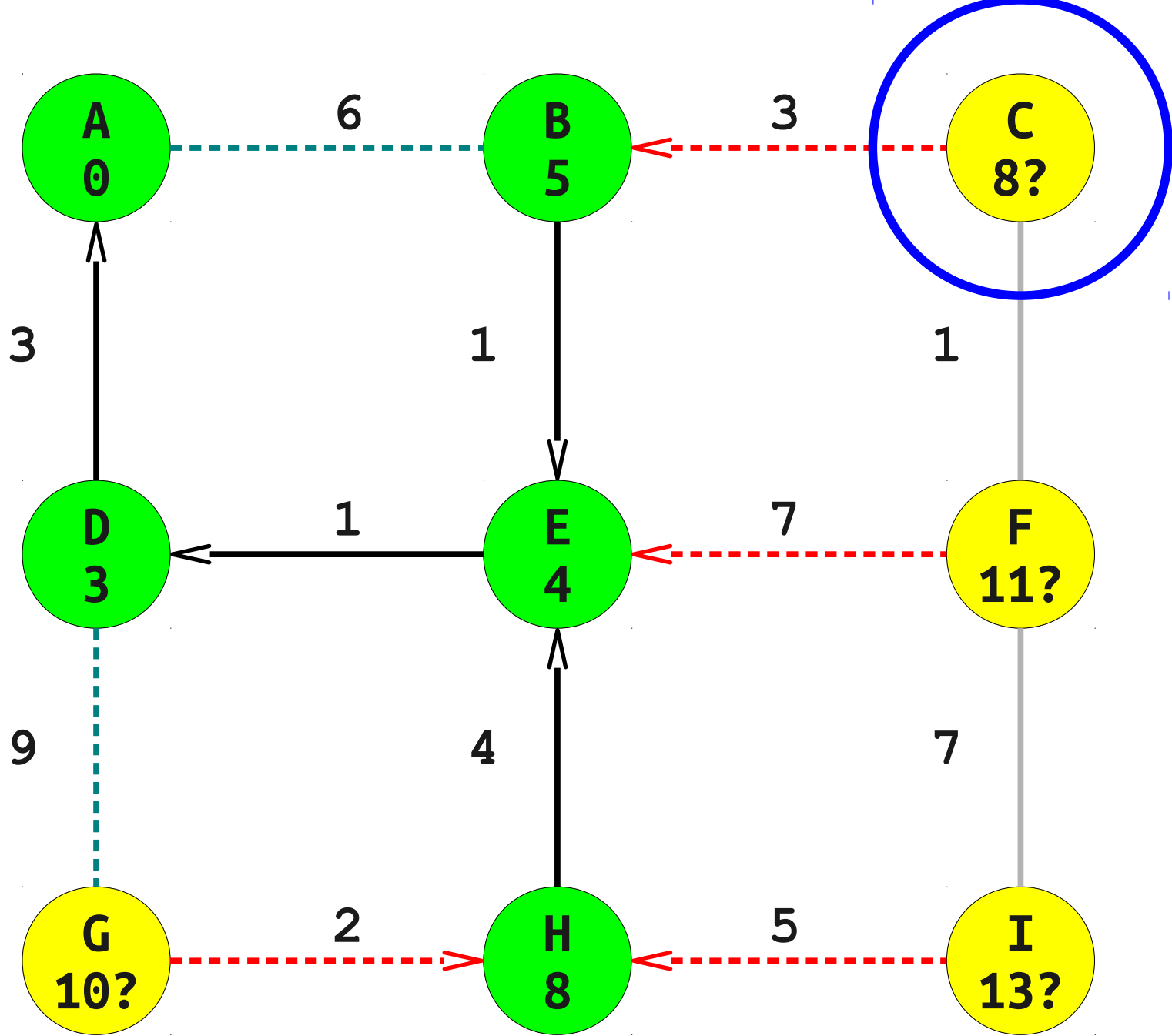
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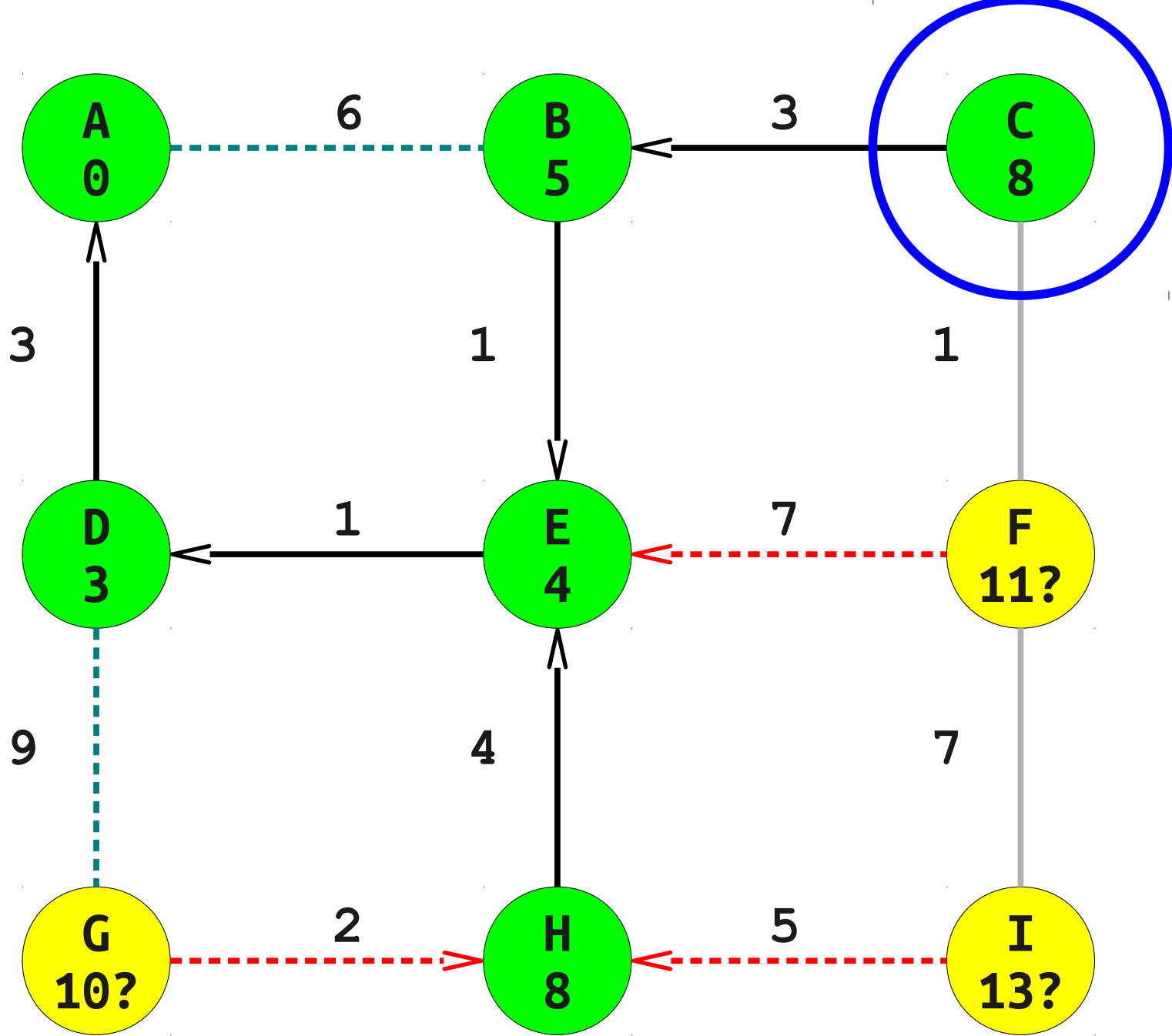
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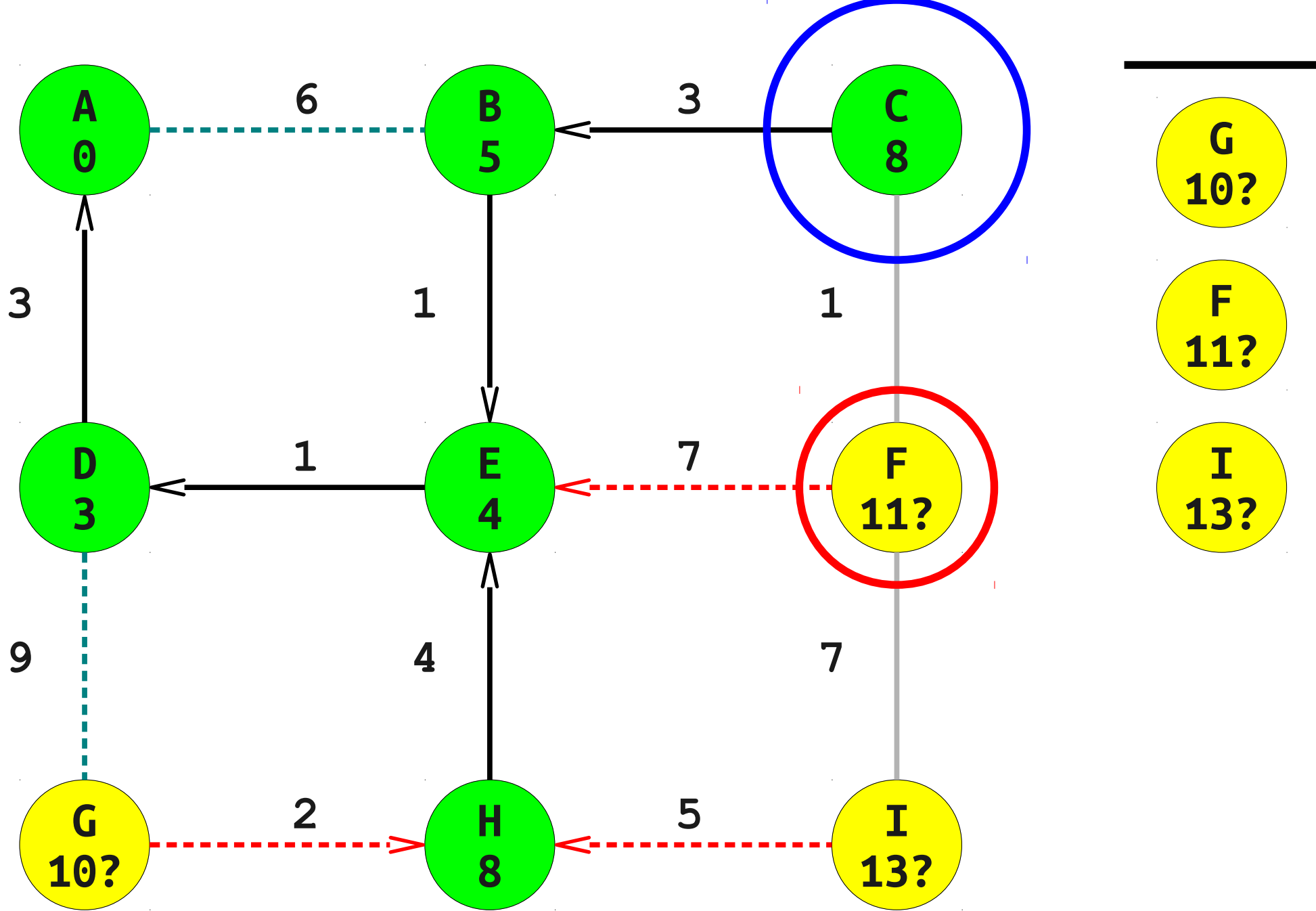
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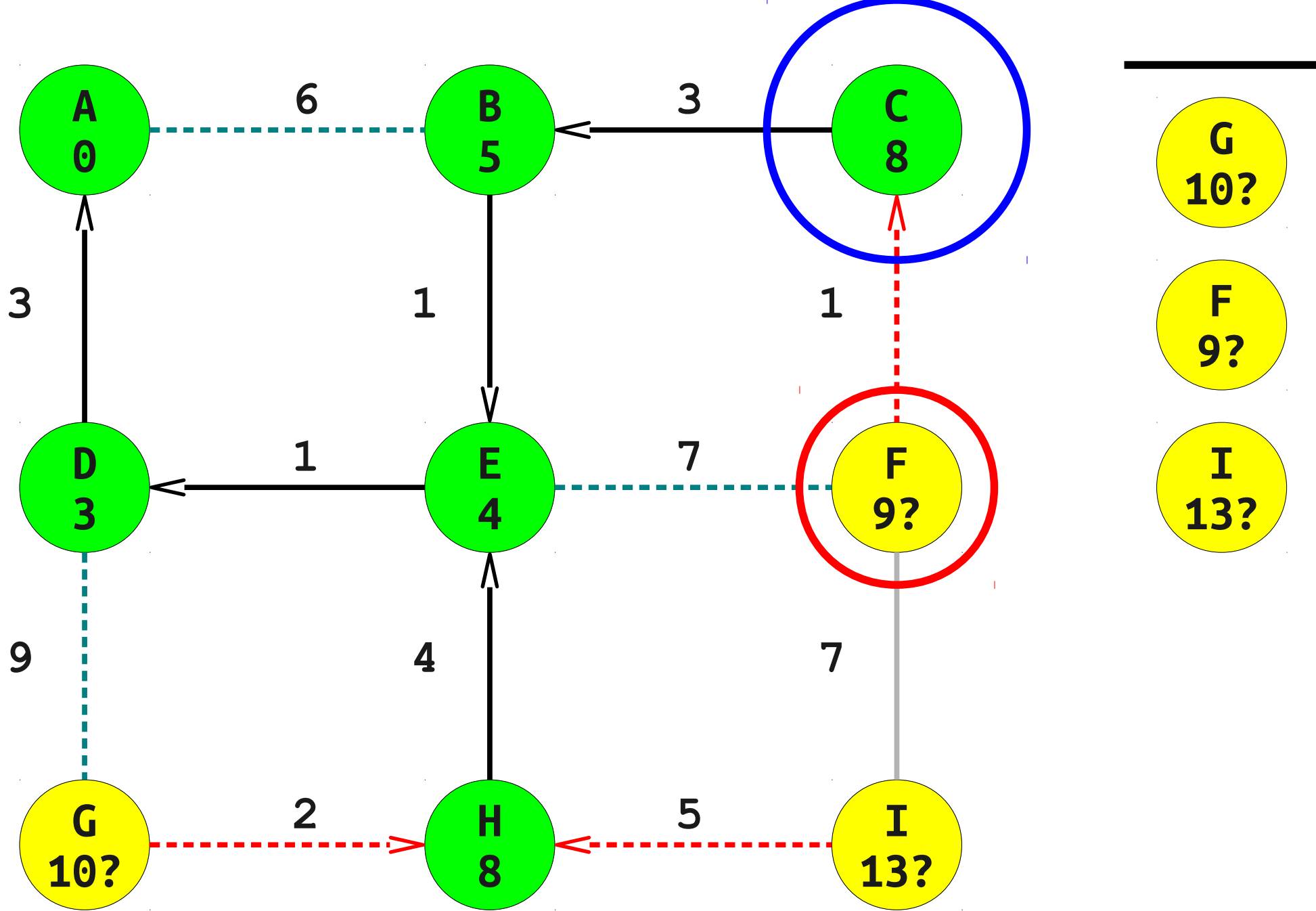
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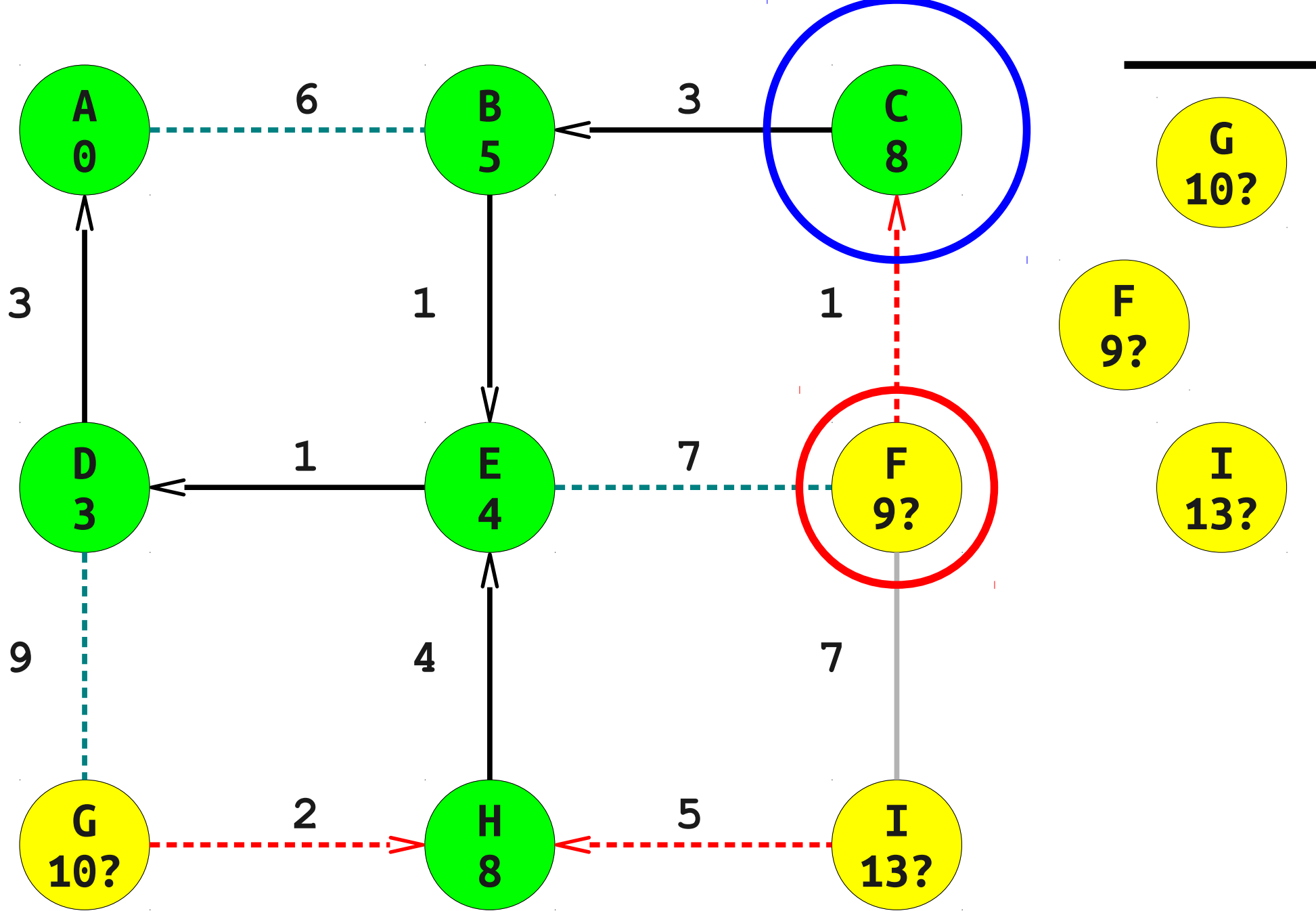
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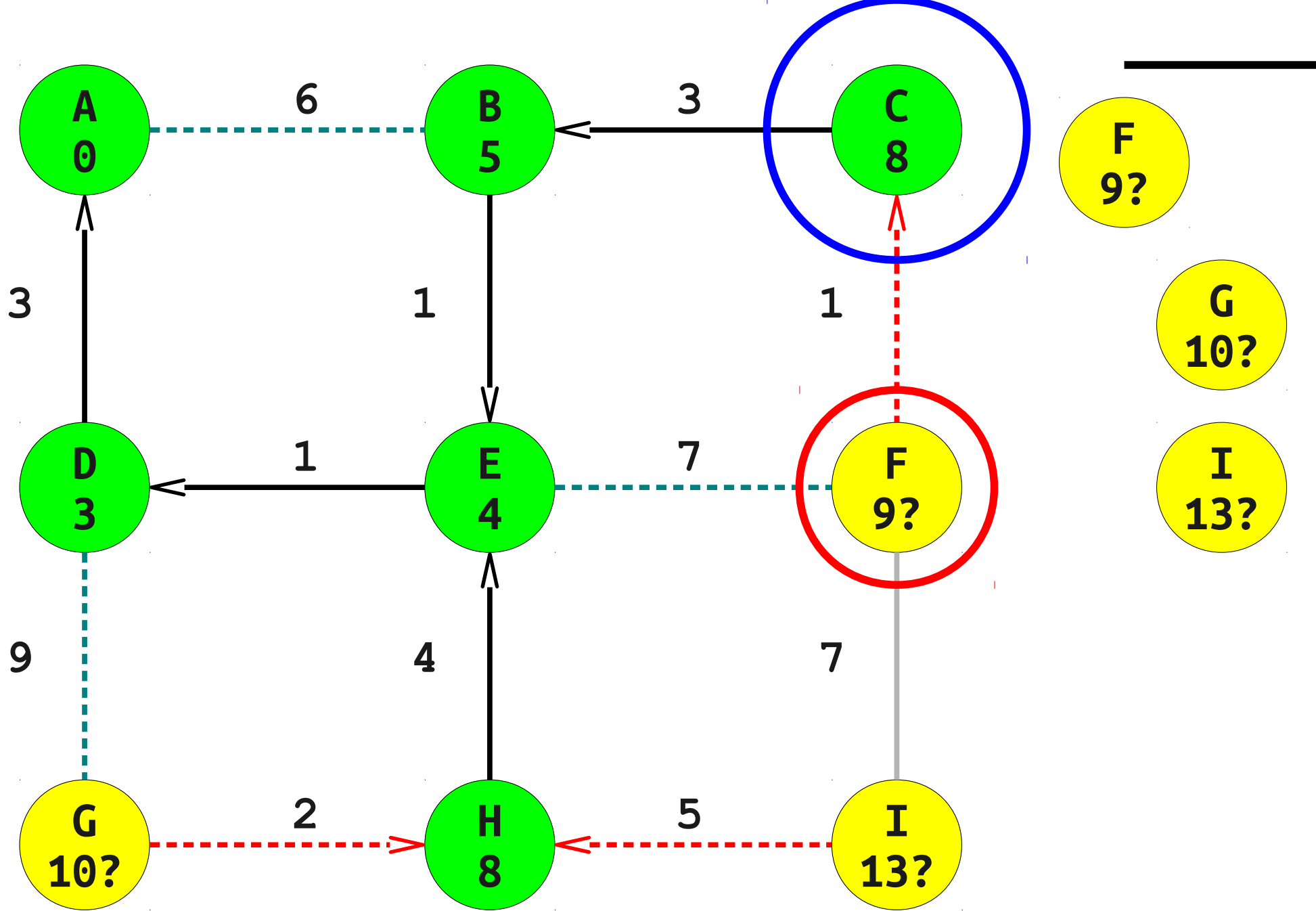


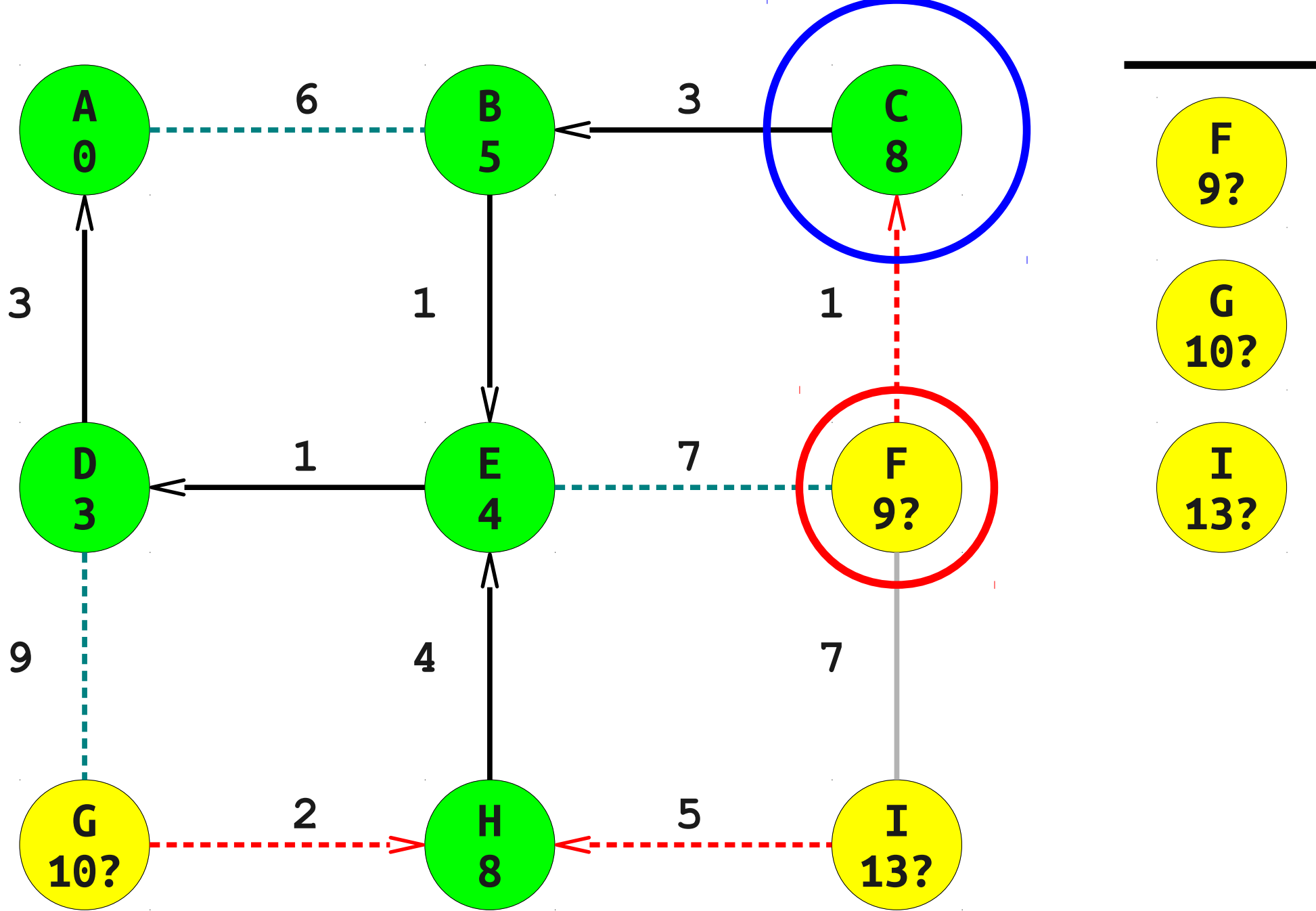
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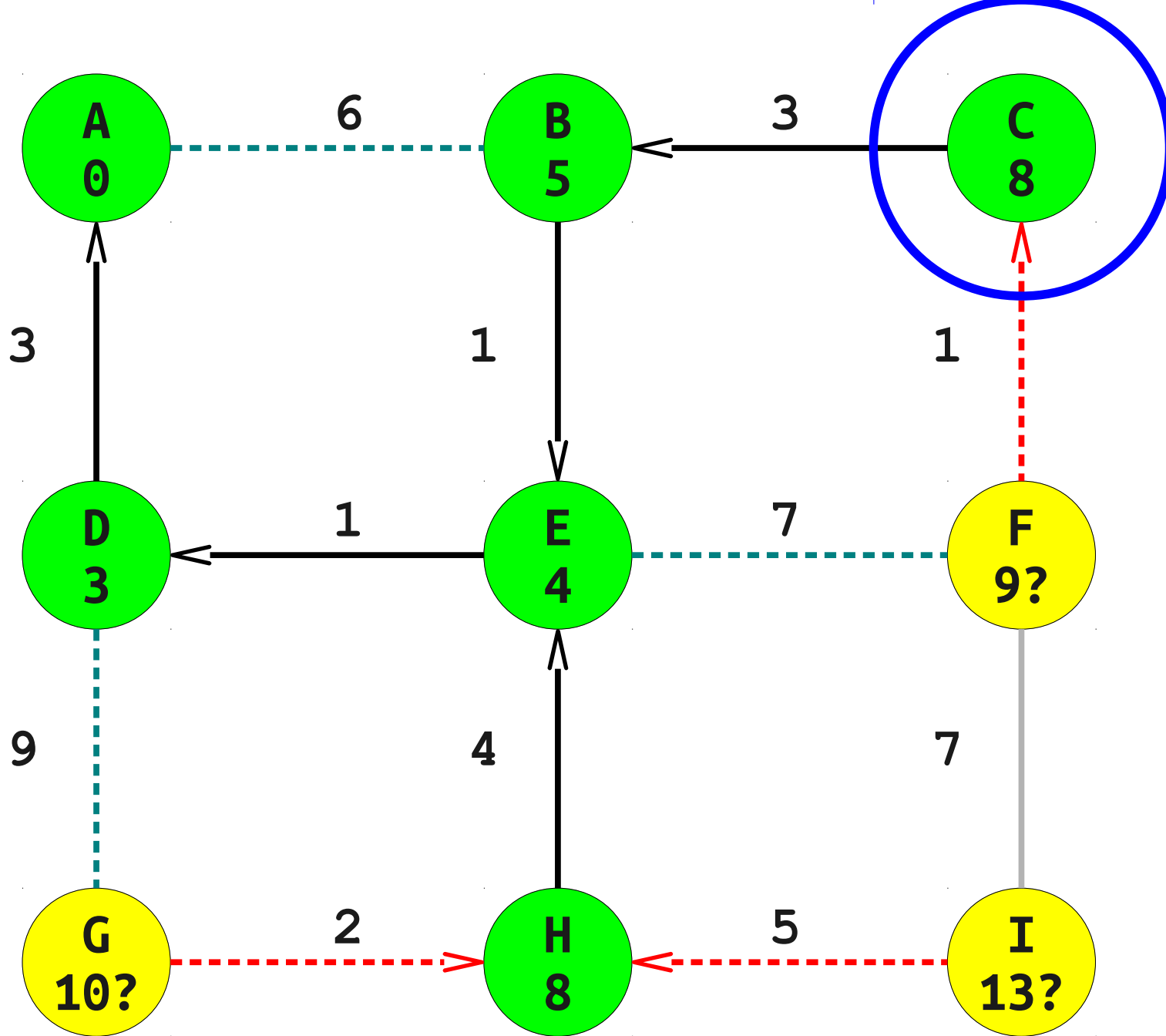
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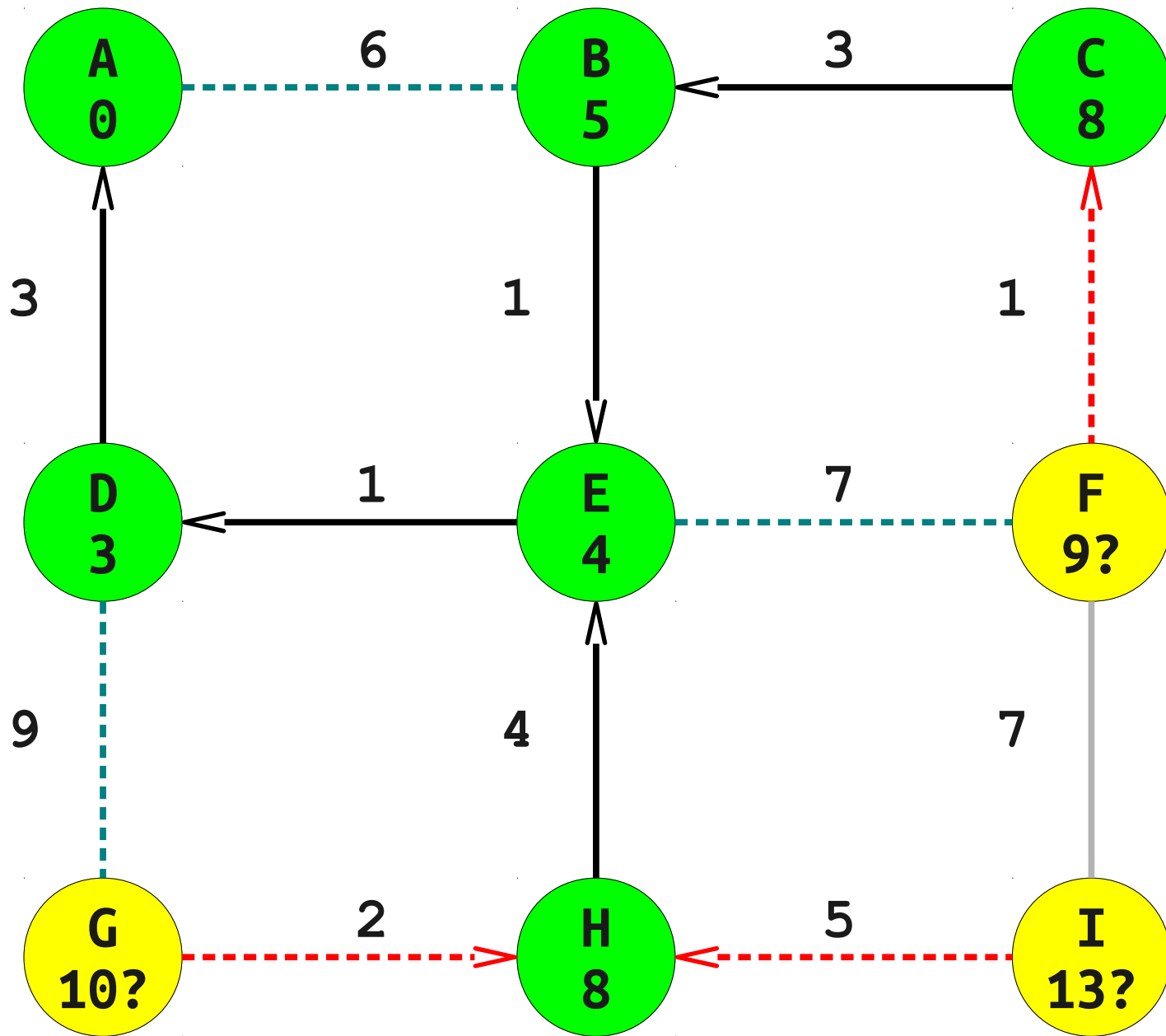




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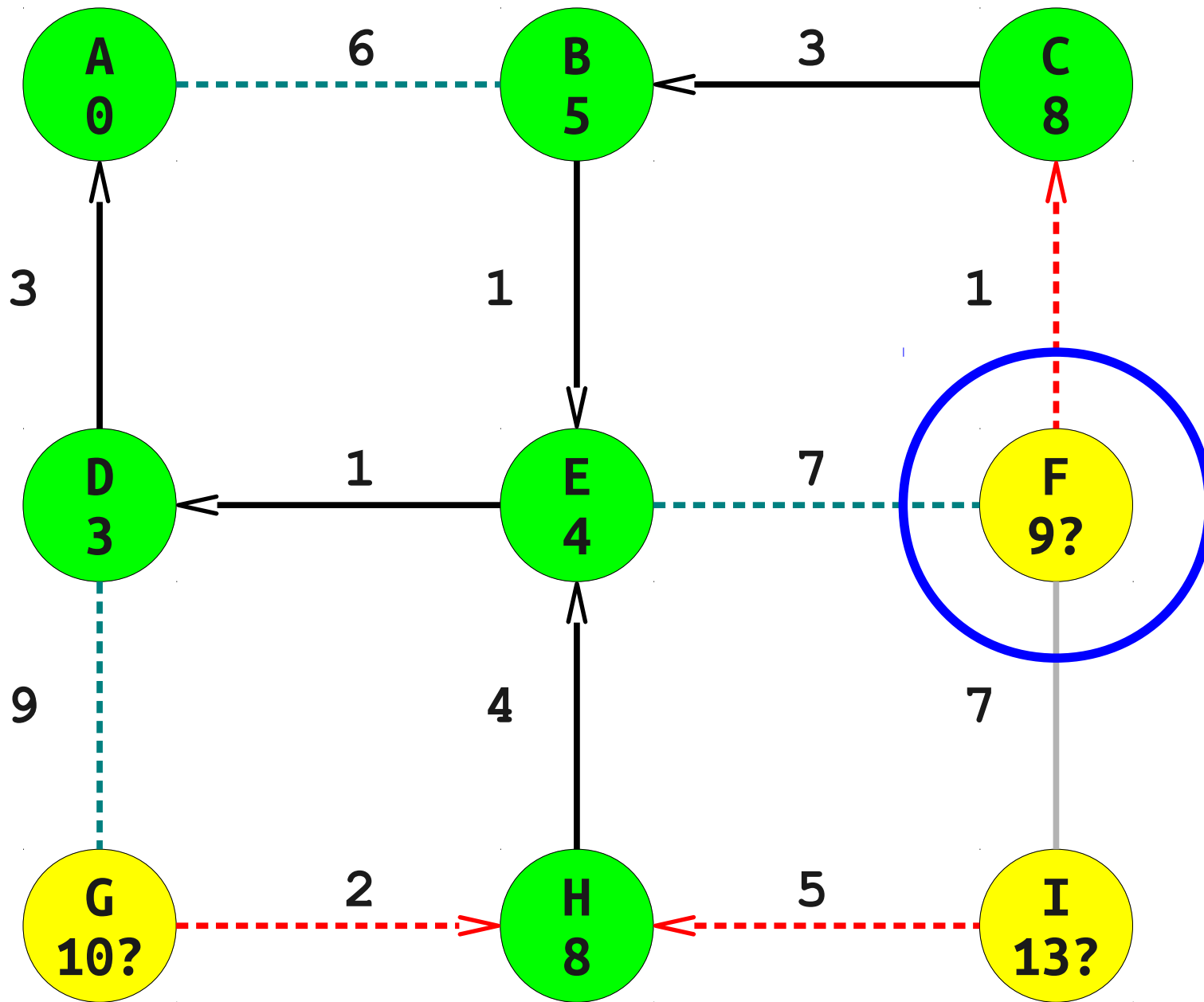
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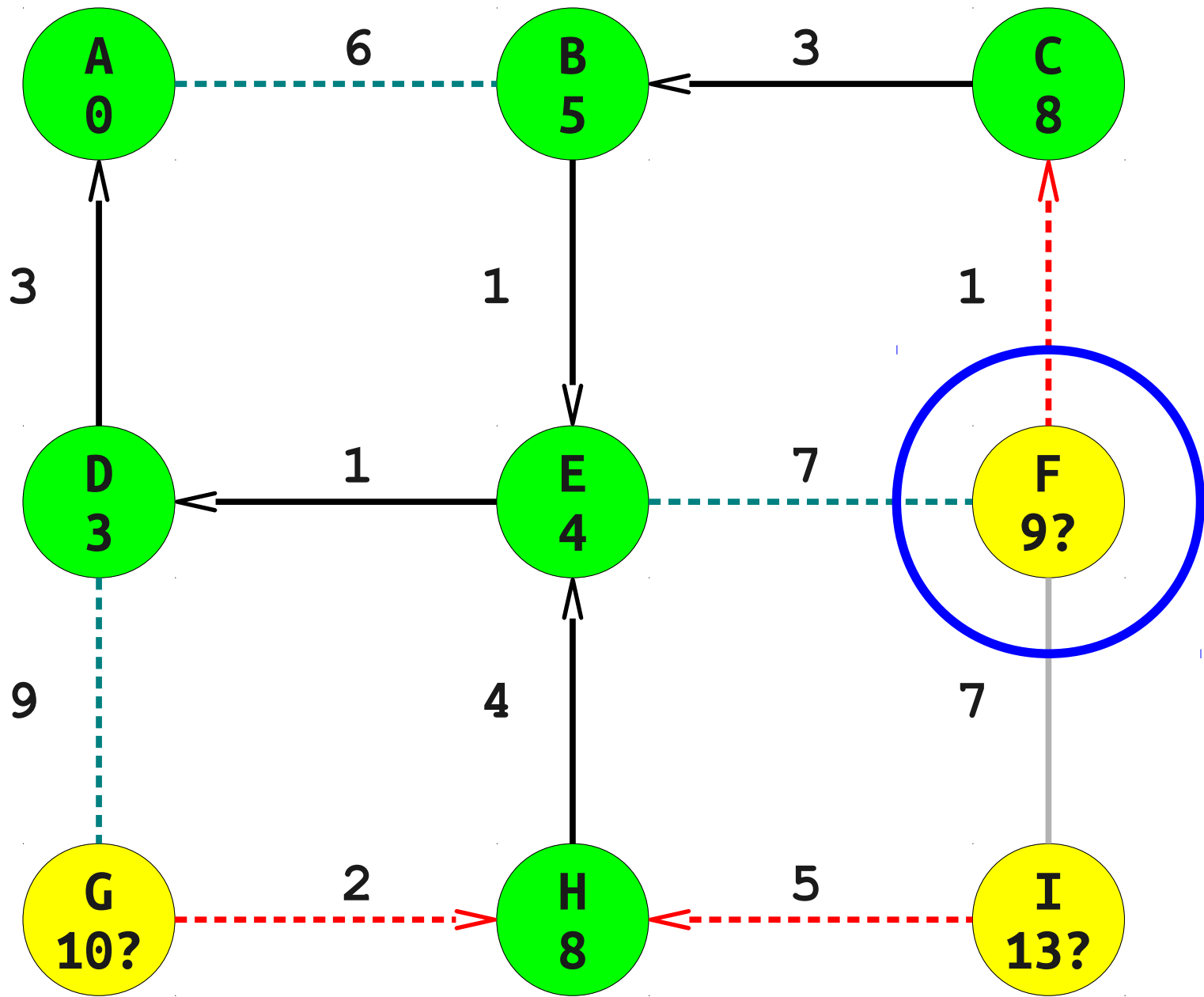
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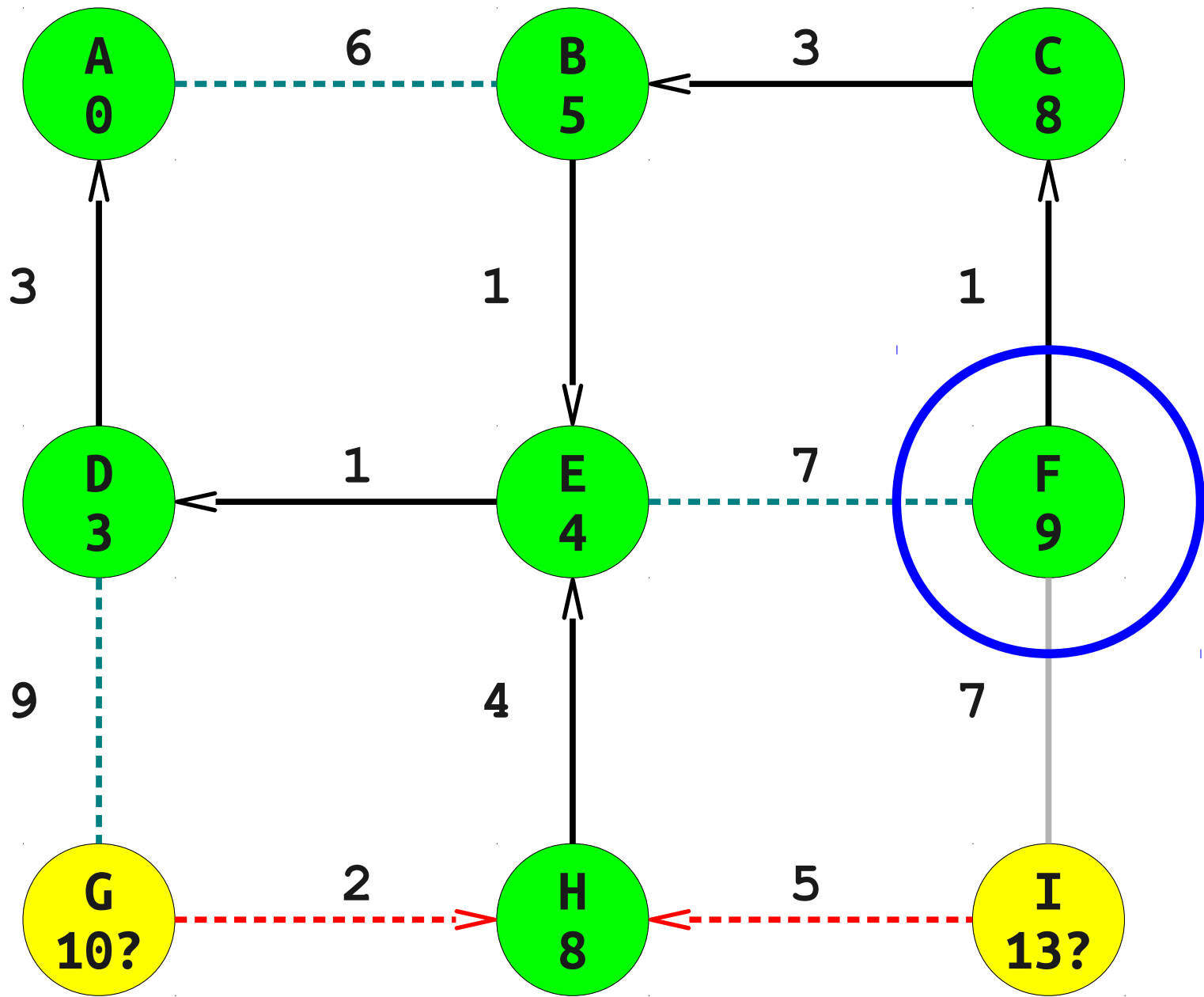
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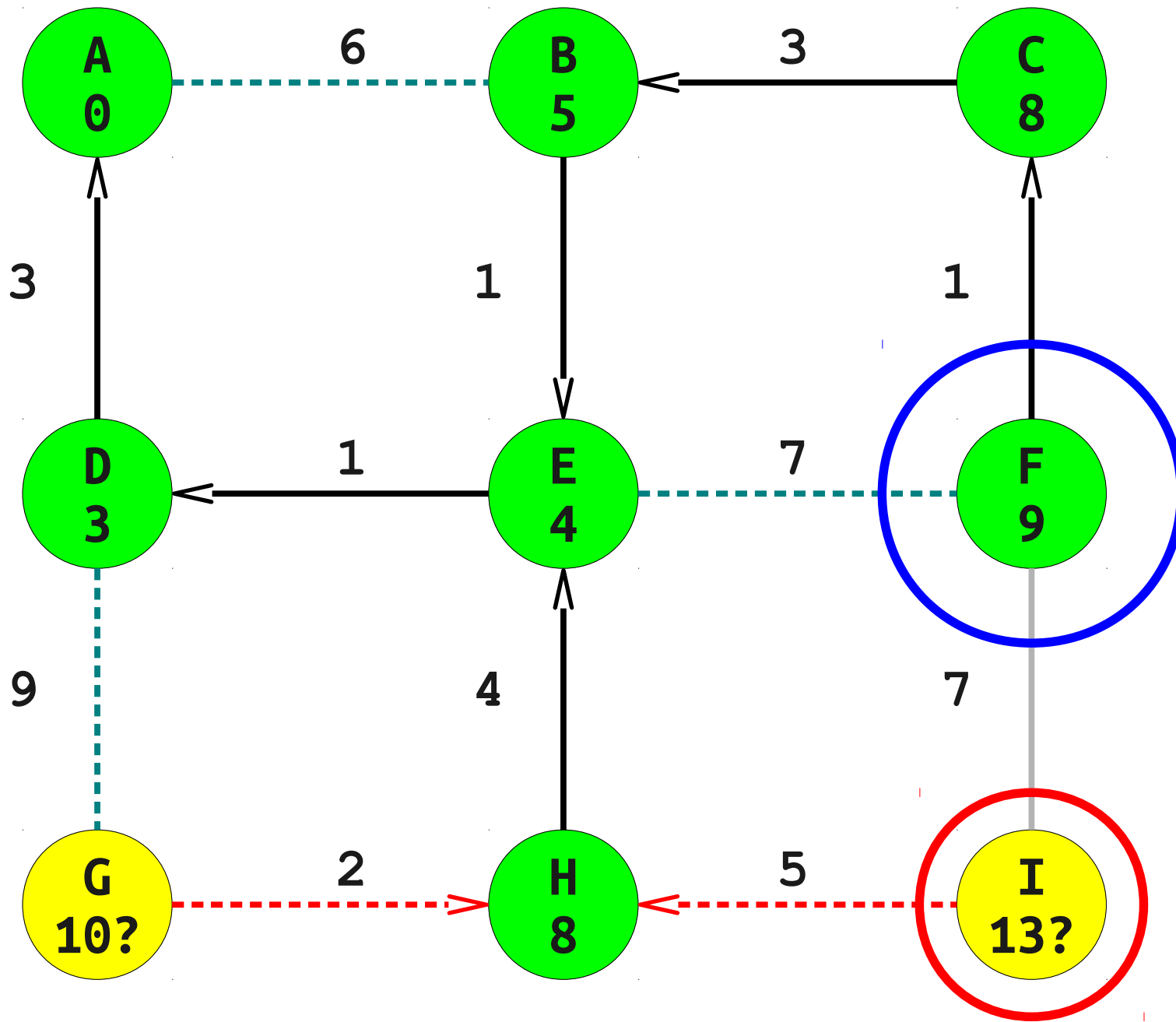
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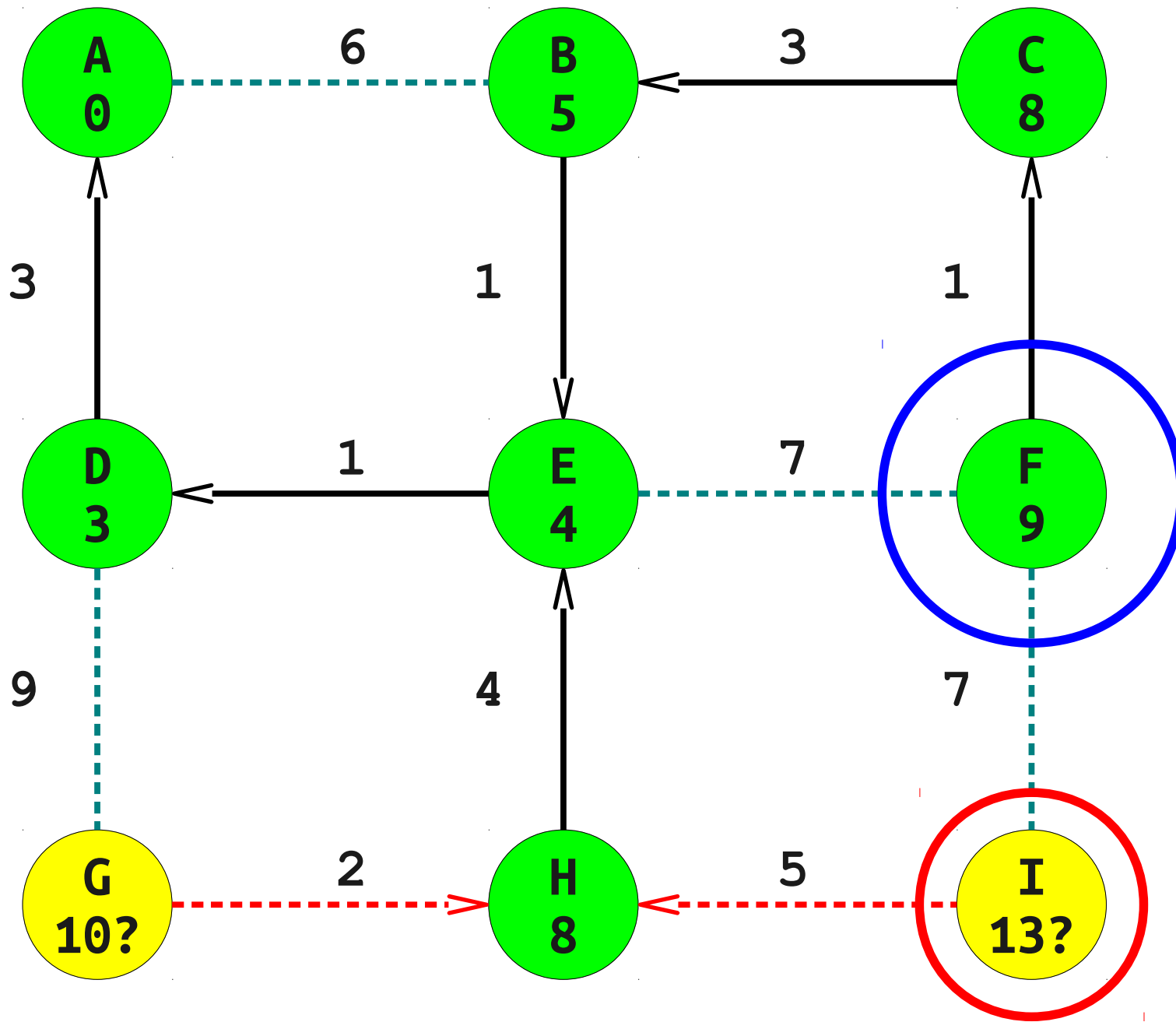
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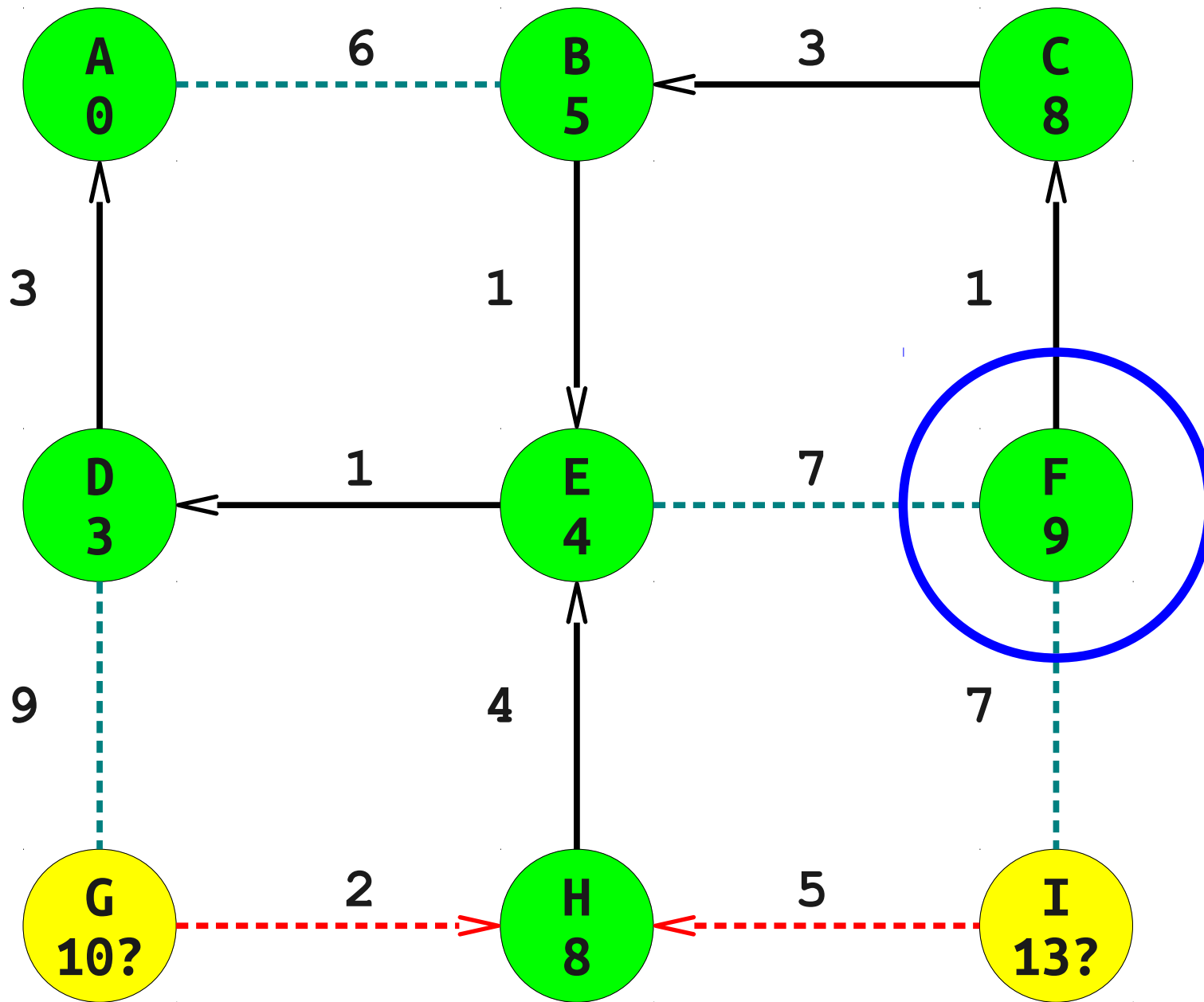
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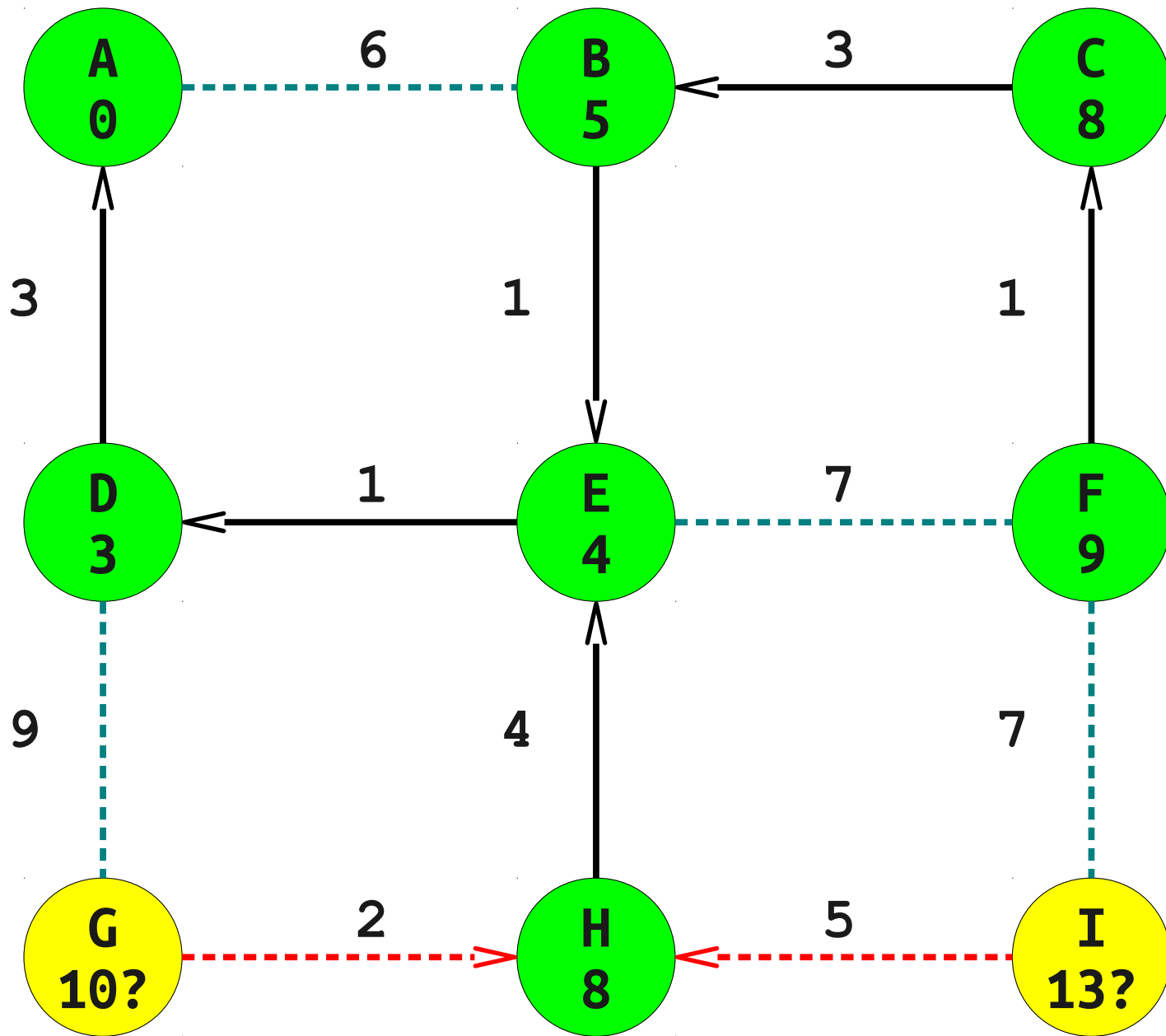
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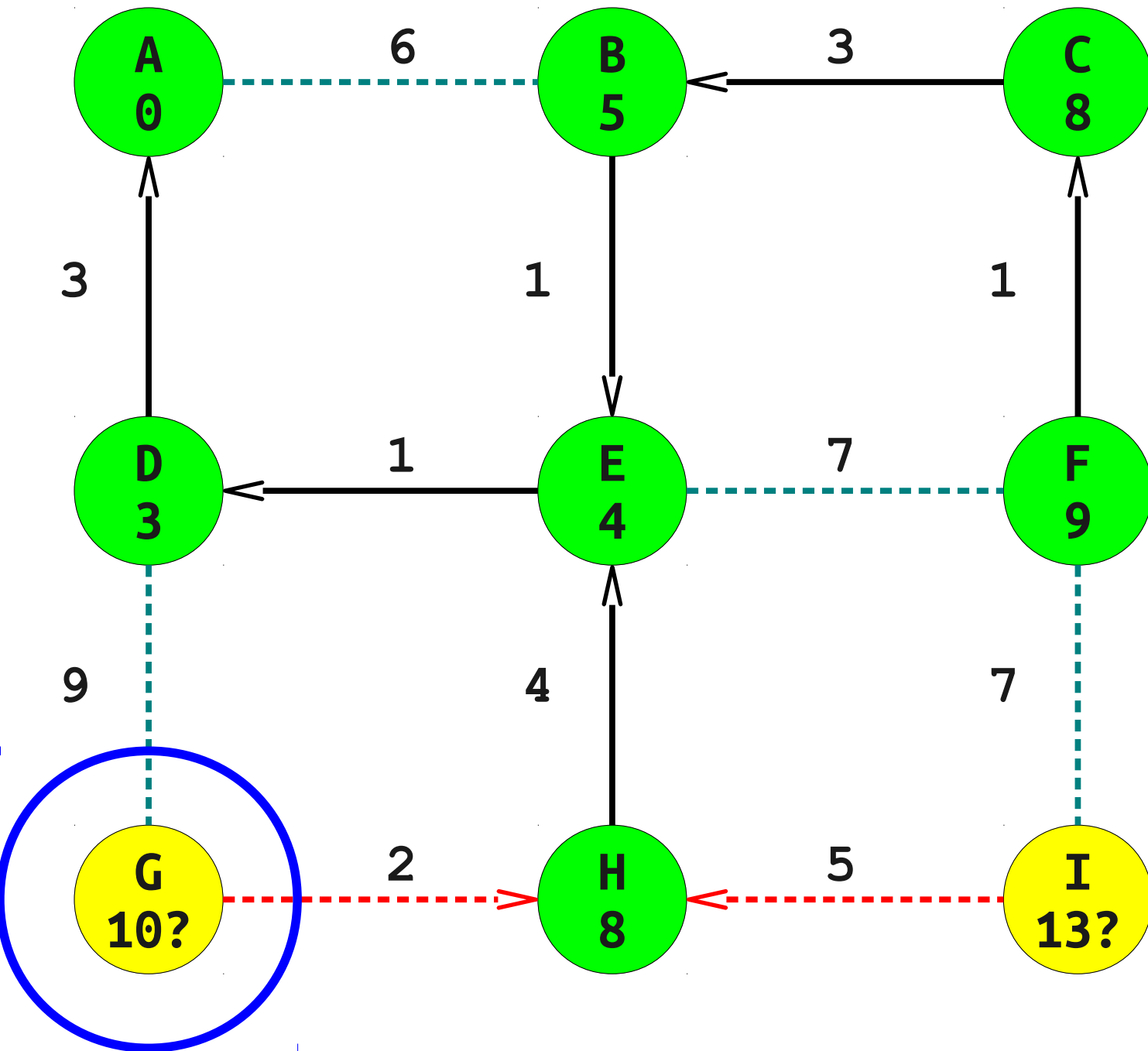
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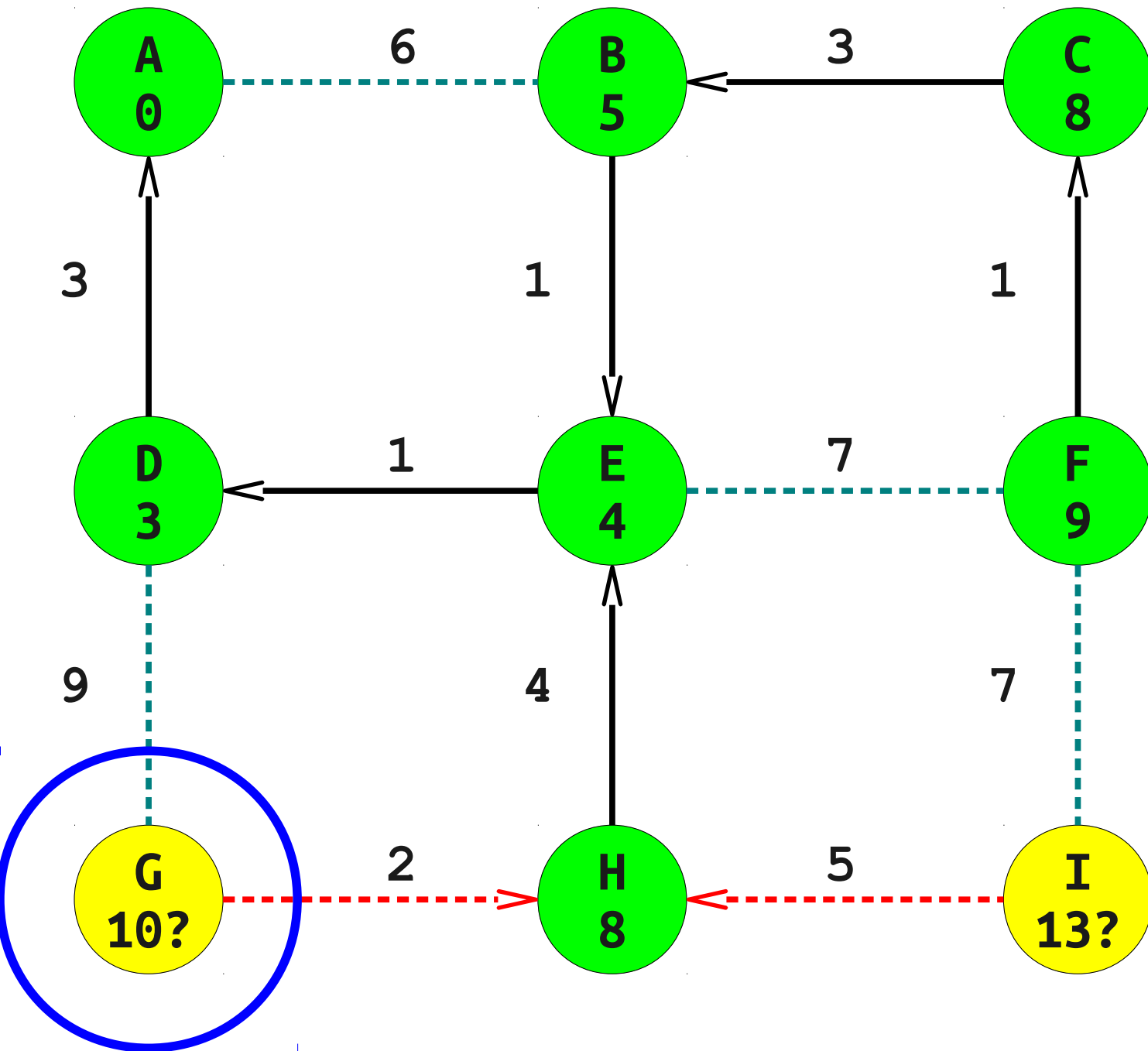
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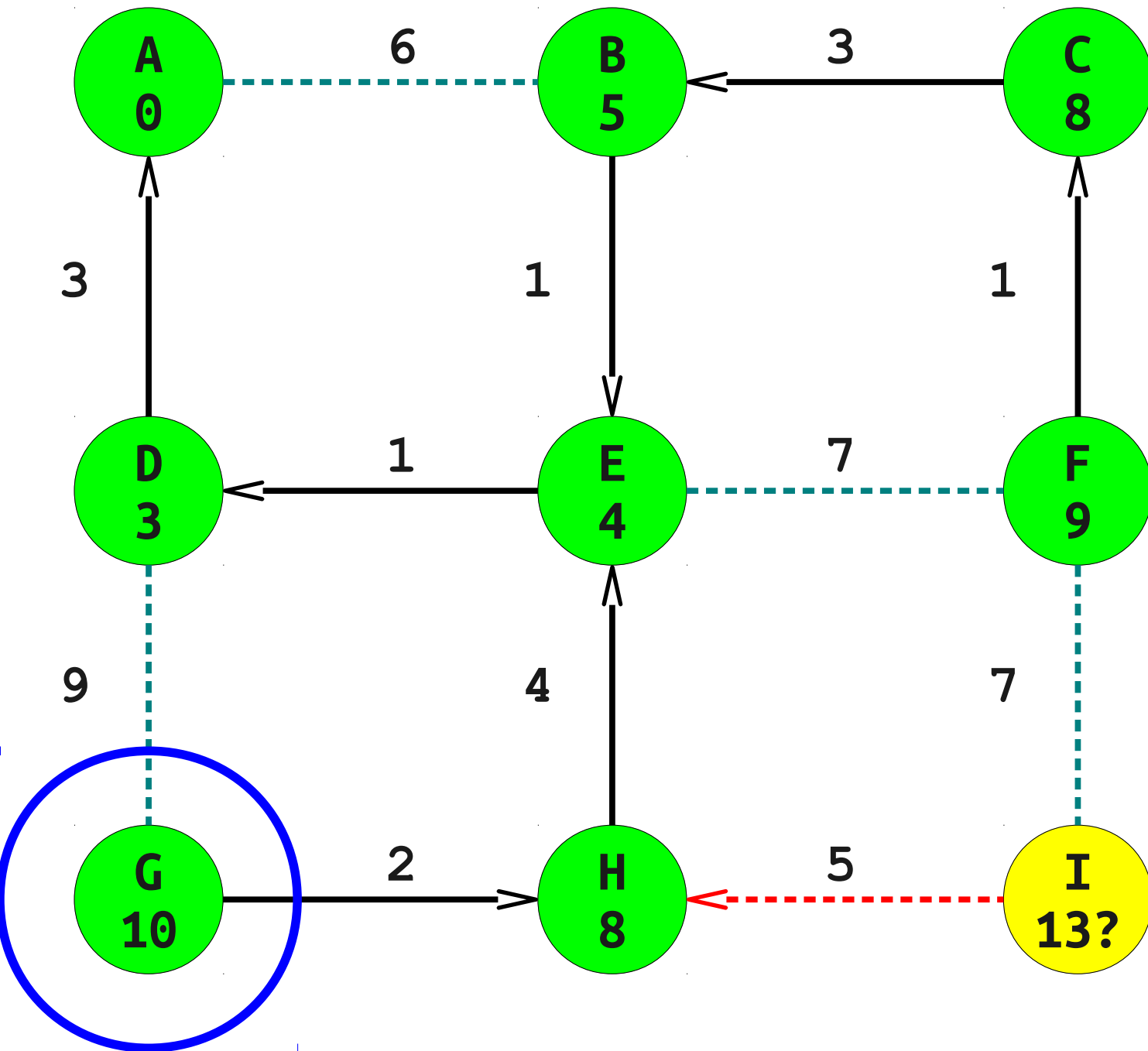


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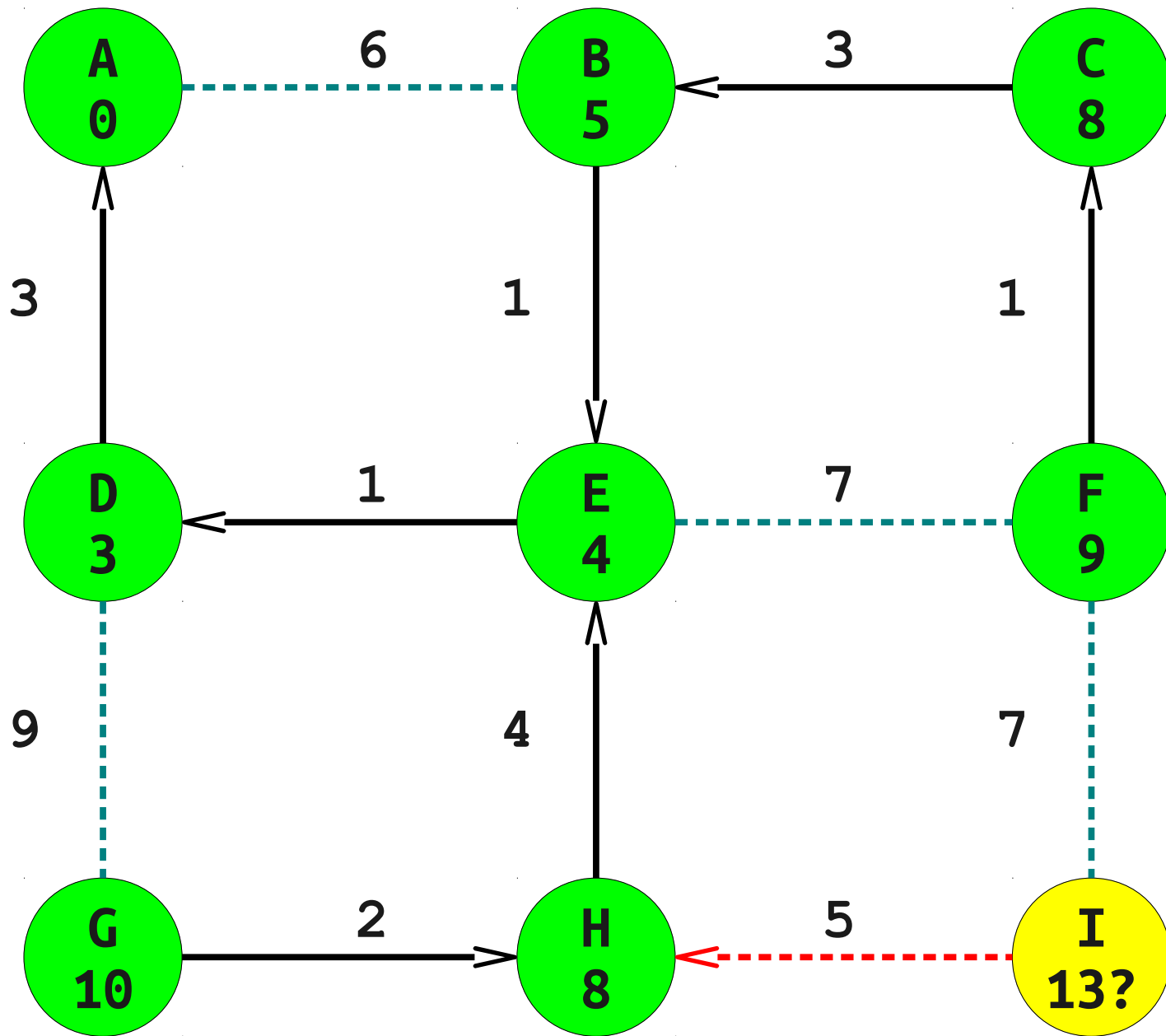
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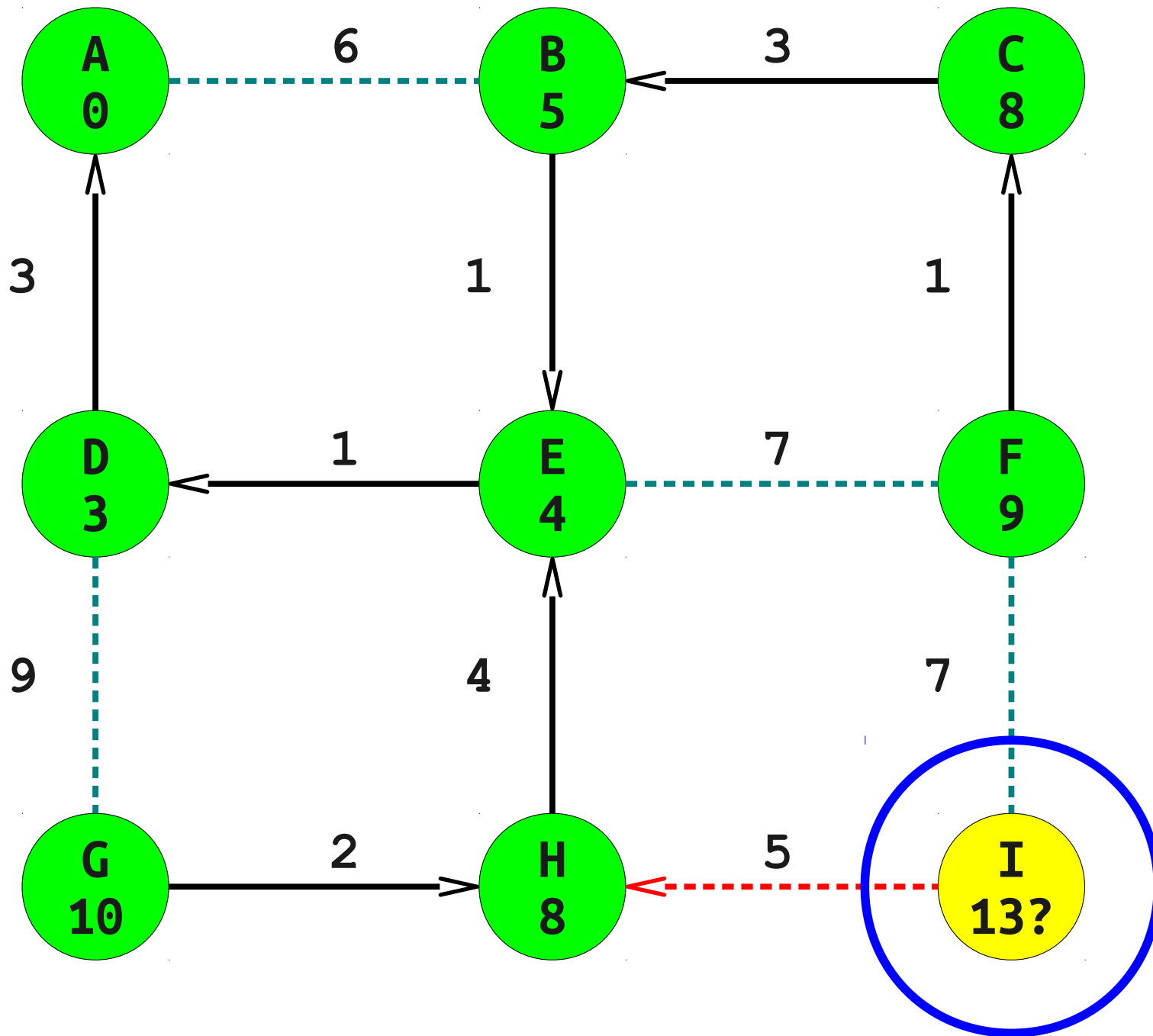
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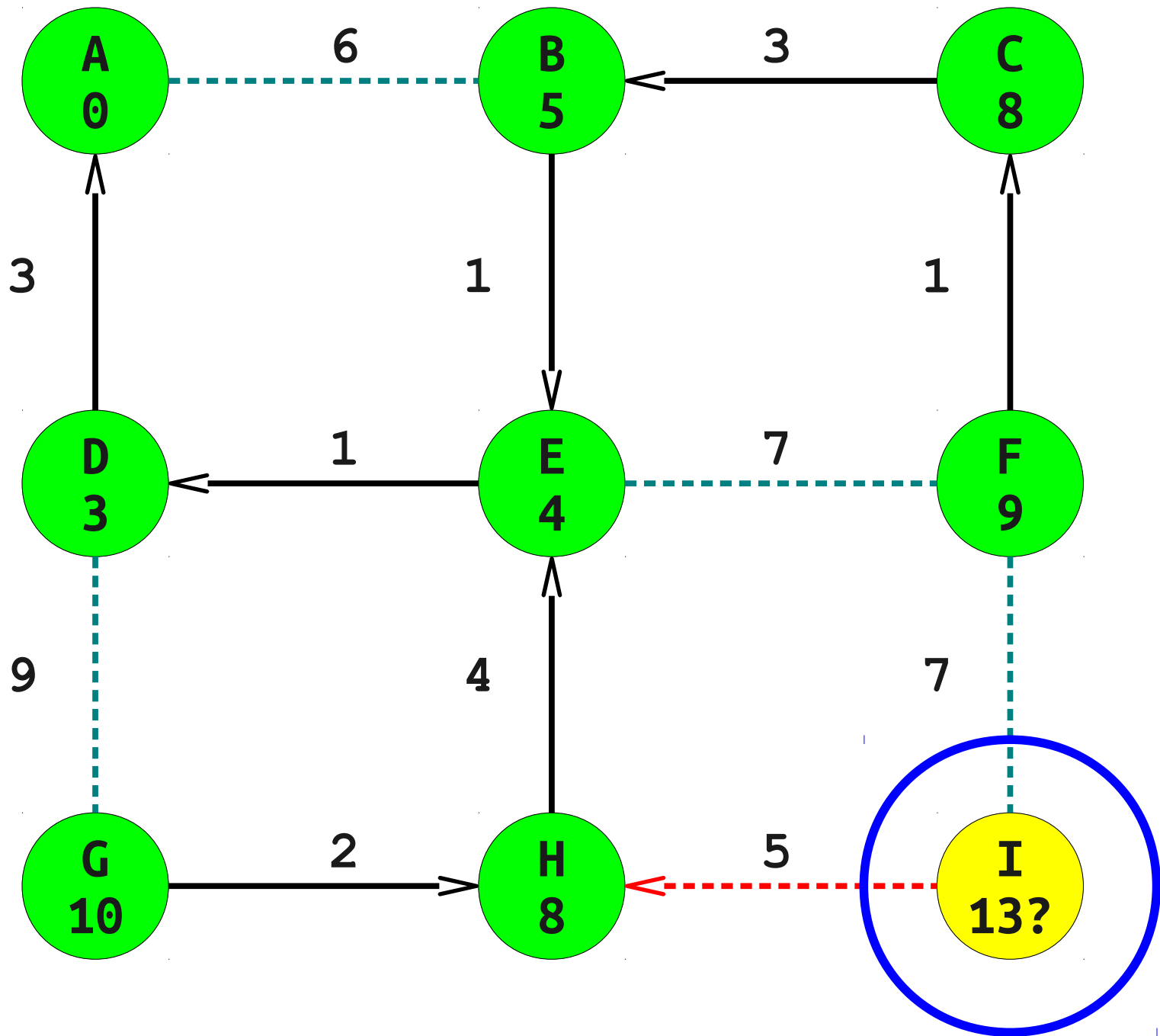
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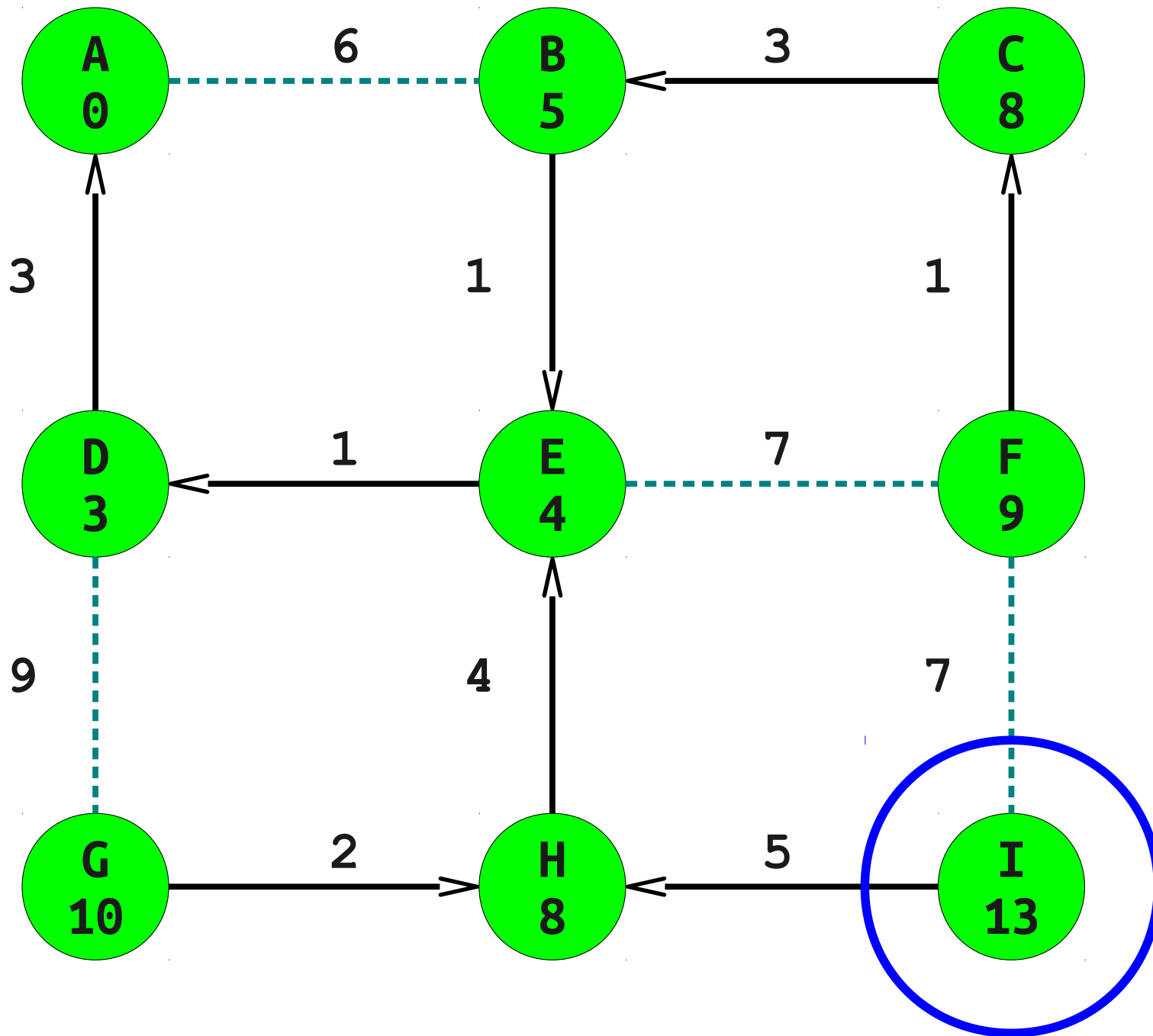


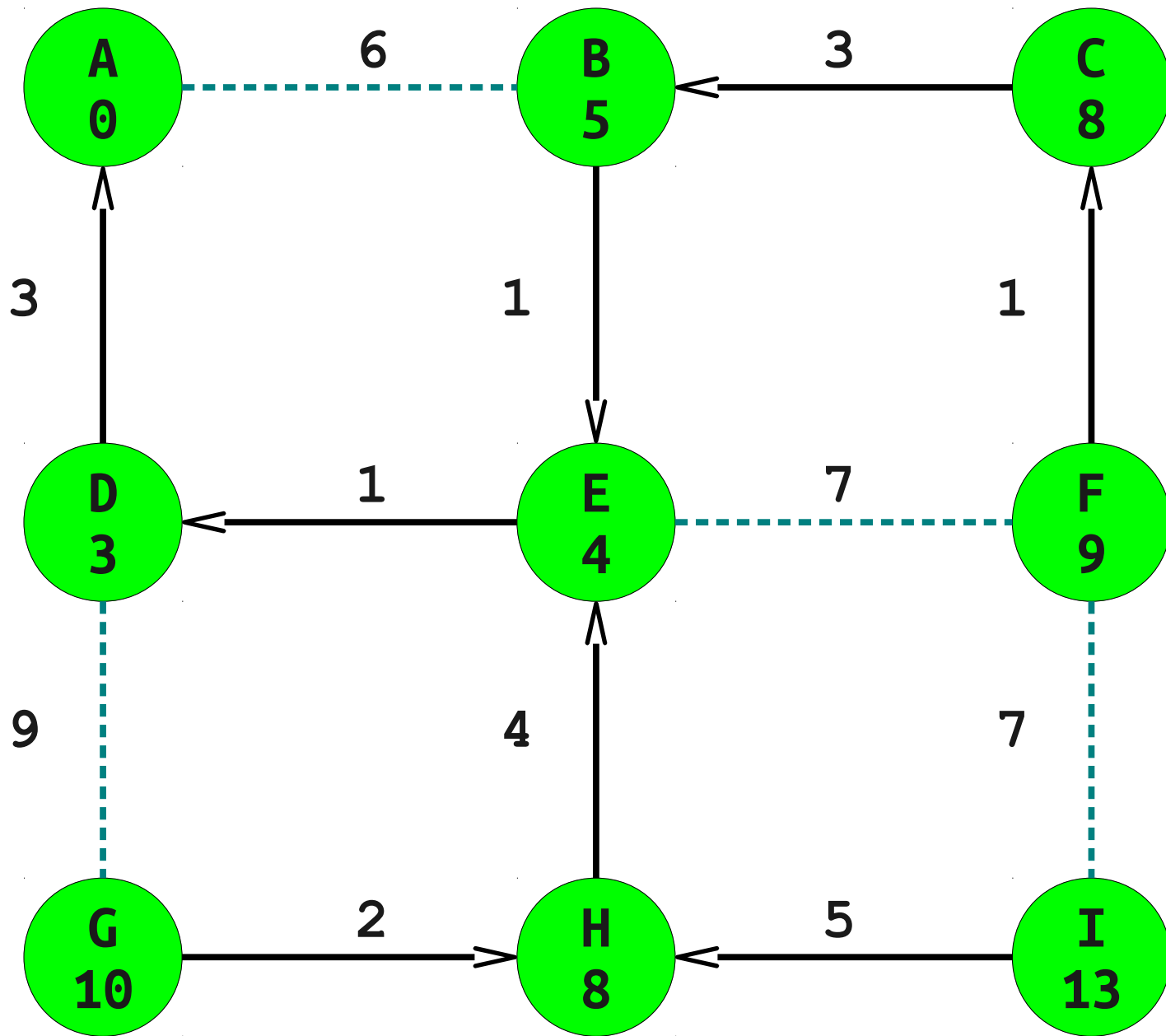
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







This approach is called
Dijkstra's algorithm.

Dijkstra's Algorithm

- Split nodes apart into three groups:
 -  Green nodes, where we already have the shortest path;
 -  Gray nodes, which we have never seen; and
 -  Yellow nodes that we still need to process.
- Dijkstra's algorithm works as follows:
 - Mark all nodes gray except the start node, which is yellow and has cost 0.
 - Until no yellow nodes remain:
 - Choose the yellow node with the lowest total cost.
 - Mark that node green.
 - Mark all its gray neighbors yellow and with the appropriate cost.
 - Update the costs of all adjacent yellow nodes by considering the path through the current node.

An Important Note

- The version of Dijkstra's algorithm I have just described is *not* the same as the version described in the course reader.
- This version is more complex than the book's version, but is much faster.

Differences from BFS

- BFS uses a queue, while Dijkstra's algorithm uses a **priority queue**, where priorities are potential distances to nodes.
 - Need a special operation called **decrease-key**, which lowers the priority of an enqueued element.
- BFS never changes a parent pointer once it is set, while Dijkstra's algorithm might change parent pointers.
 - If a possible path to a node is found to be incorrect, then its parent might change.

The next slide will be really useful as a reference, so I changed the color scheme so that you can find it more easily.

- Mark all nodes as gray.
- Mark the initial node **s** as yellow and at candidate distance **0**.
- Enqueue **s** into the priority queue.
- While not all nodes have been visited:
 - Dequeue the lowest-cost node **u** from the priority queue.
 - Color **u** green. The candidate distance **d** that is currently stored for node **u** is the length of the shortest path from **s** to **u**.
 - If **u** is the destination node **t**, you have found the shortest path from **s** to **t** and are done.
 - For each node **v** connected to **u** by an edge of length **L**:
 - If **v** is gray:
 - Color **v** yellow.
 - Mark **v**'s distance as **d + L**.
 - Set **v**'s parent to be **u**.
 - Enqueue **v** into the priority queue.
 - If **v** is yellow and the candidate distance to **v** is greater than **d + L**:
 - Update **v**'s candidate distance to be **d + L**; the older candidate distance is incorrect.
 - Update **v**'s parent to be **u**.
 - Update **v**'s priority in the priority queue to **d + L**.

This slide is not very important, so we're
back to the old color scheme.

Implementation Concerns

- At each point in time, the algorithm must keep track of the following:
 - Which nodes are green (already done), yellow (in the priority queue but not done yet), and gray (never visited before).
 - The guessed distances to each node.
 - The parents of each yellow / green node.
 - The contents of the priority queue.
- There are *many* ways that you can store this information. It depends on many factors, including how the graph is represented.

A Few Quick Announcements

Friday Four Square!

Last one of the quarter! Today at 4:15PM
at Gates Computer Science

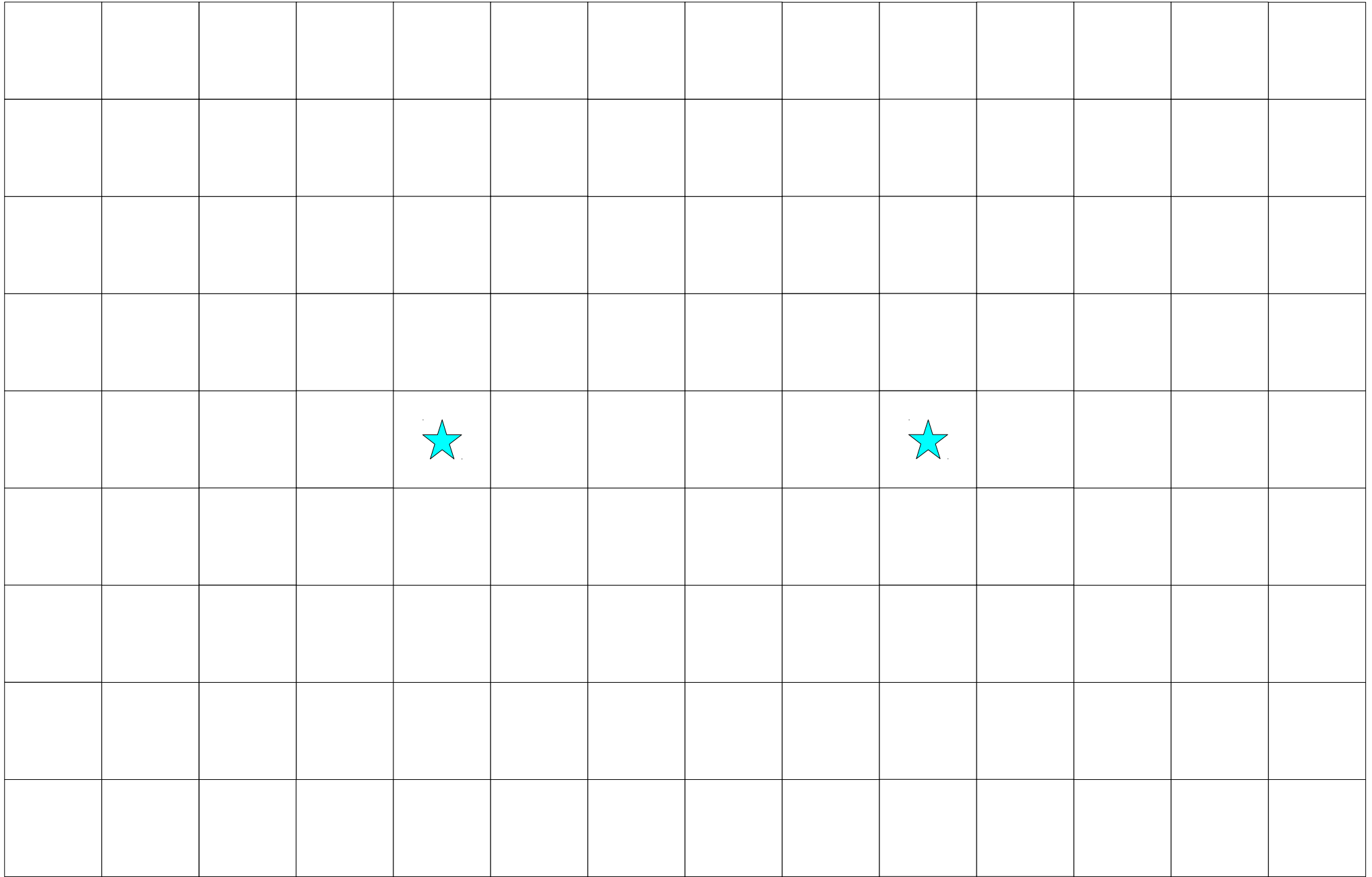
Please evaluate this course on Axess.

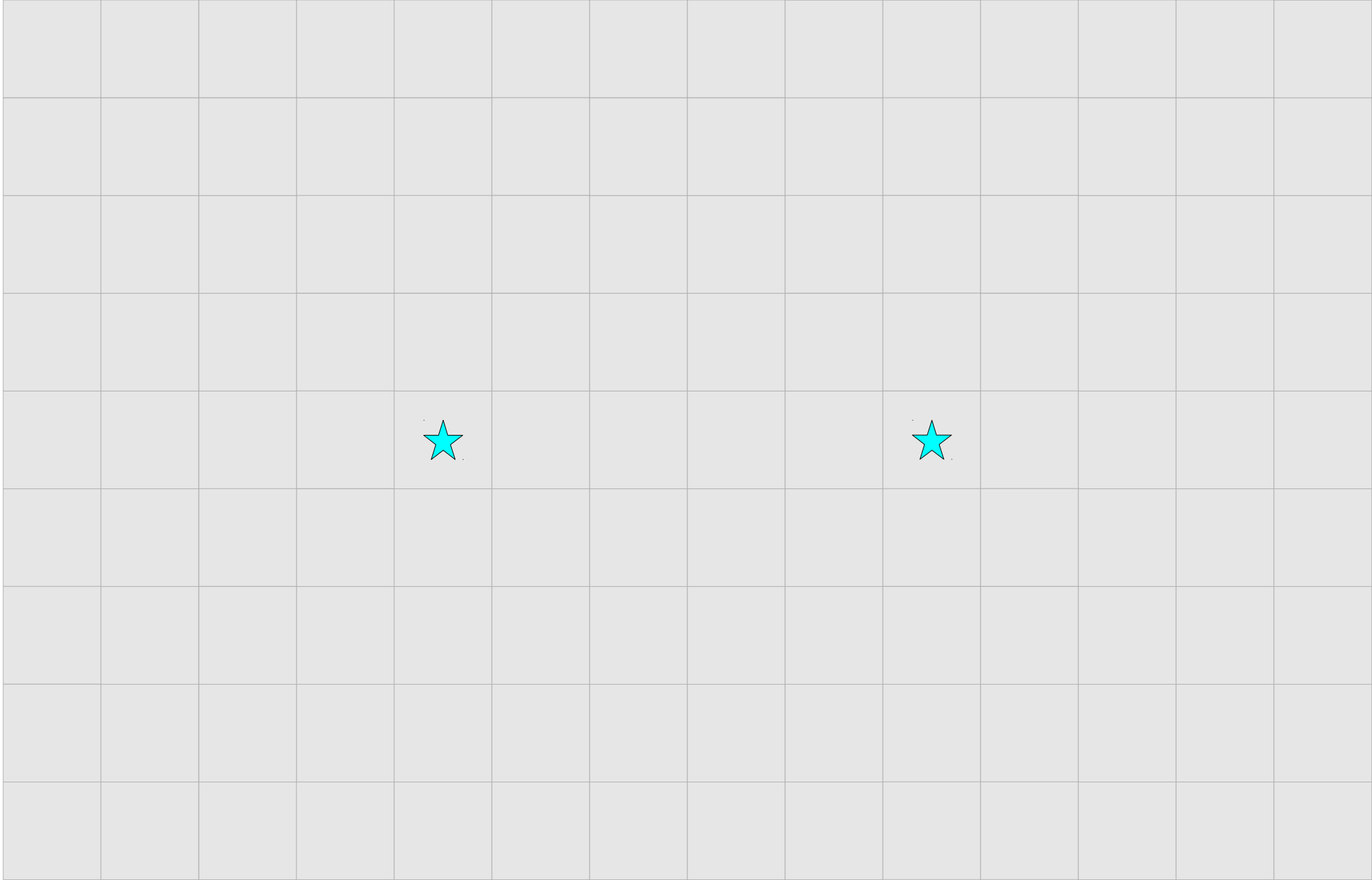
Your feedback really does make a difference. I read every single evaluation. Everything is done anonymously.

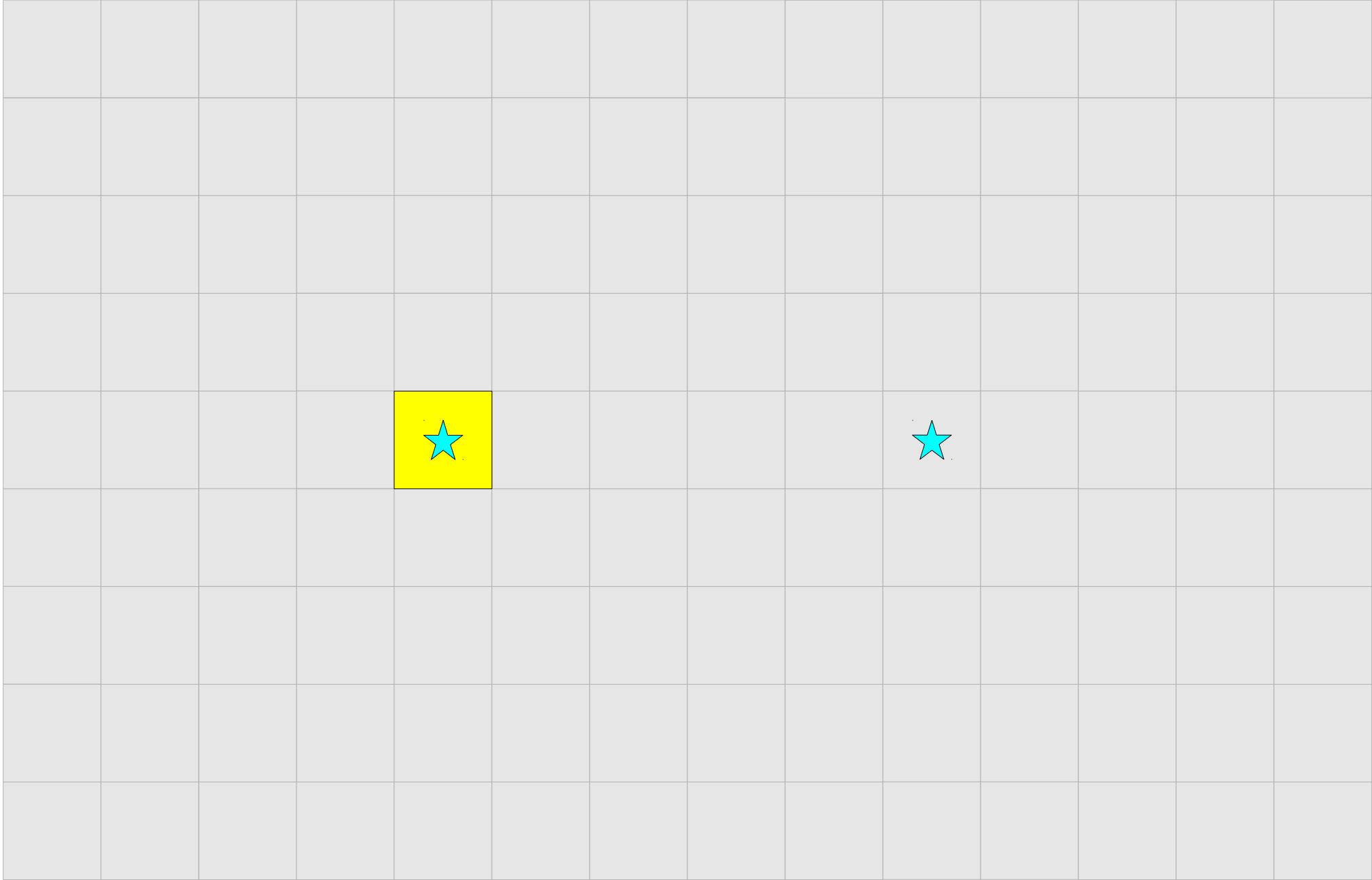
Heads-Up: Assignment 7

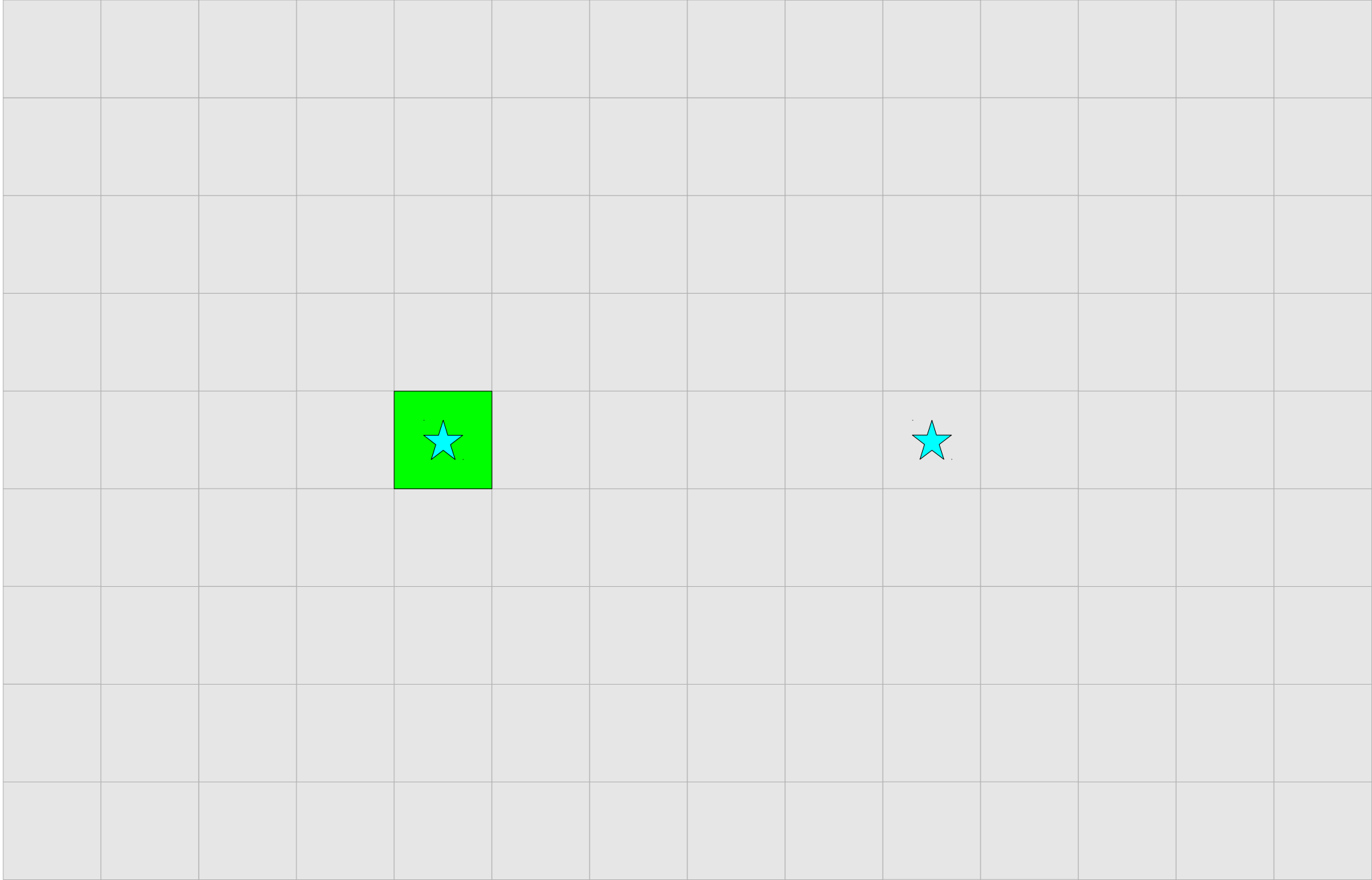
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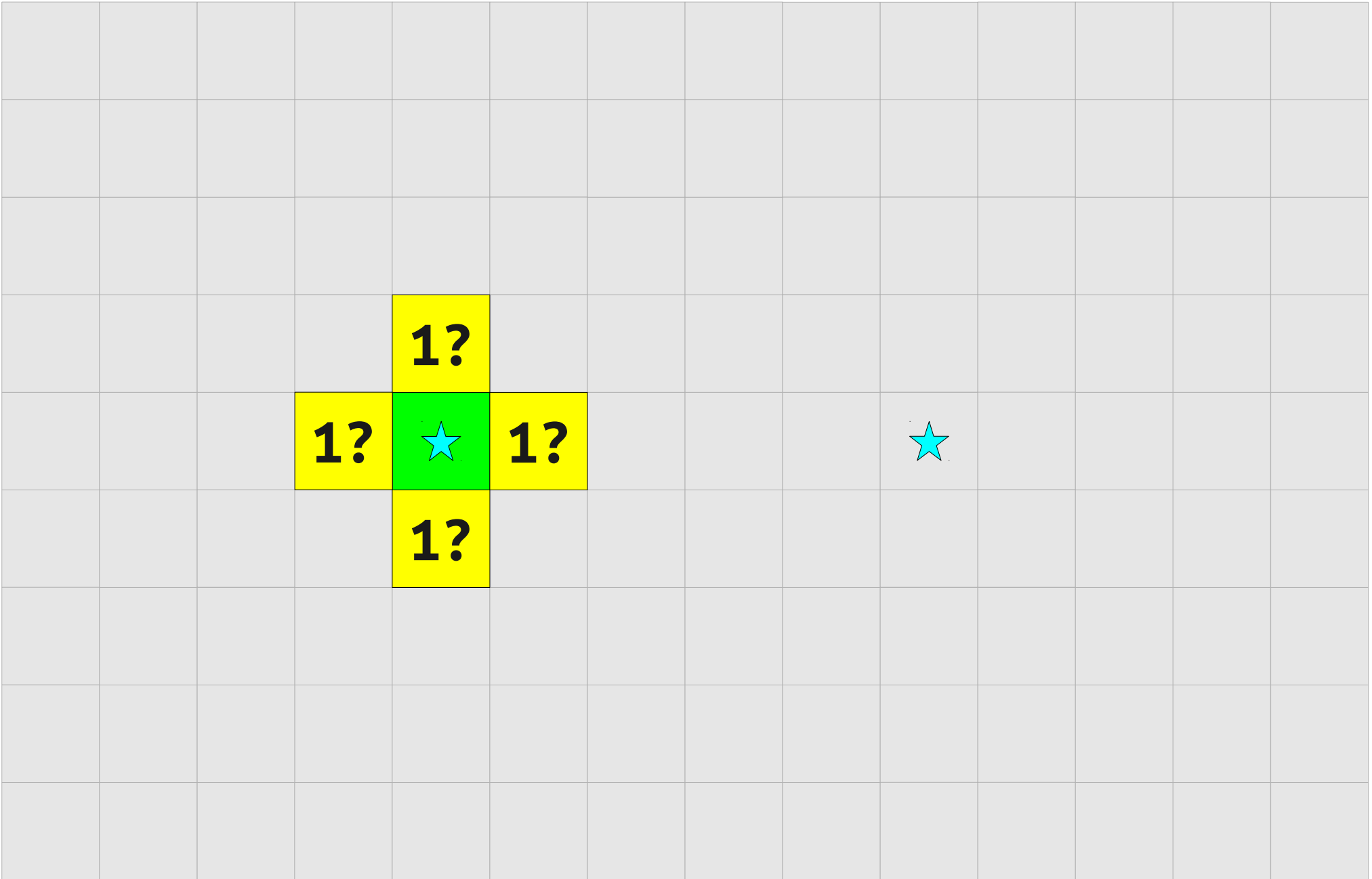
One Detail with Dijkstra's Algorithm





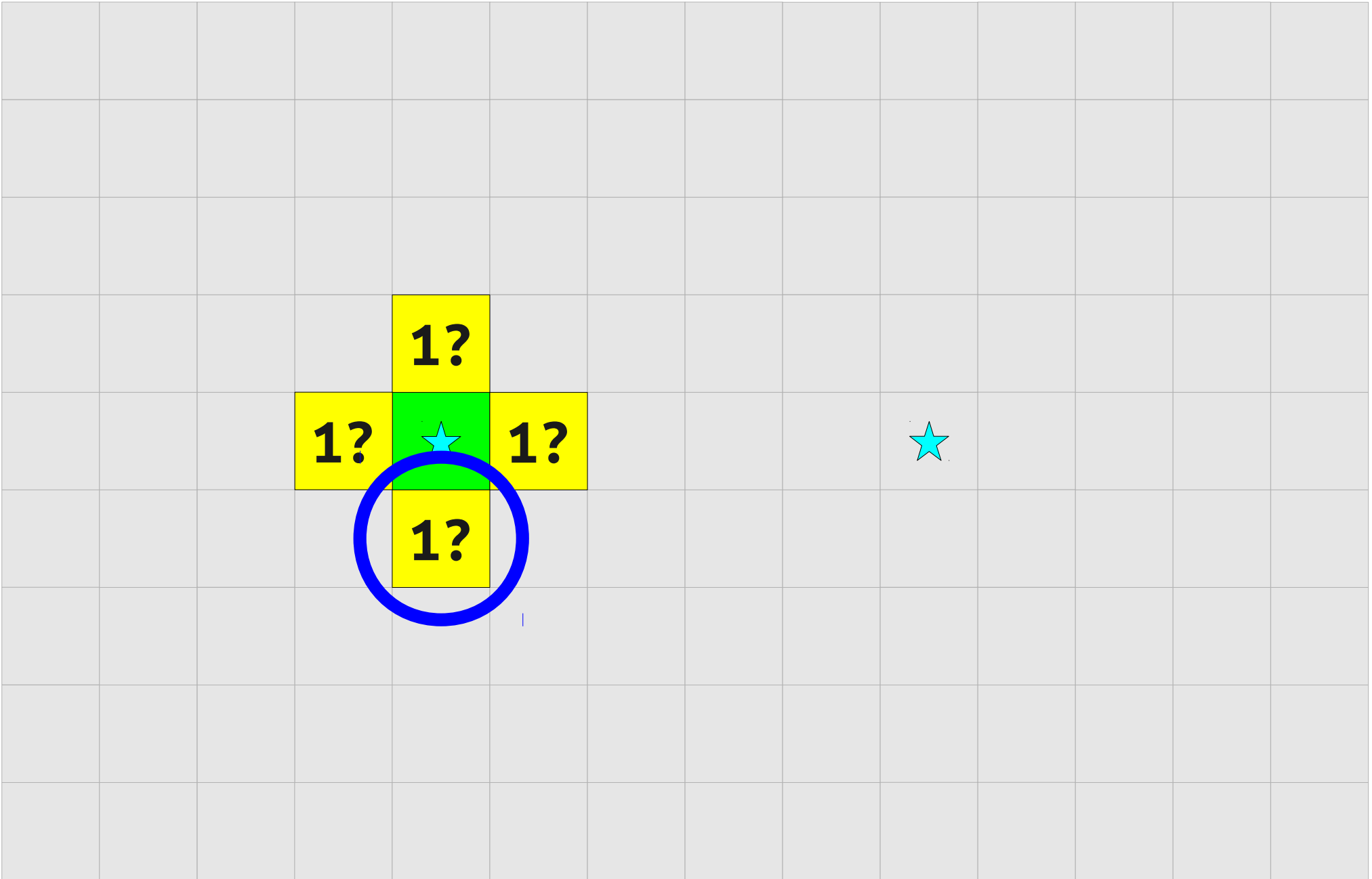


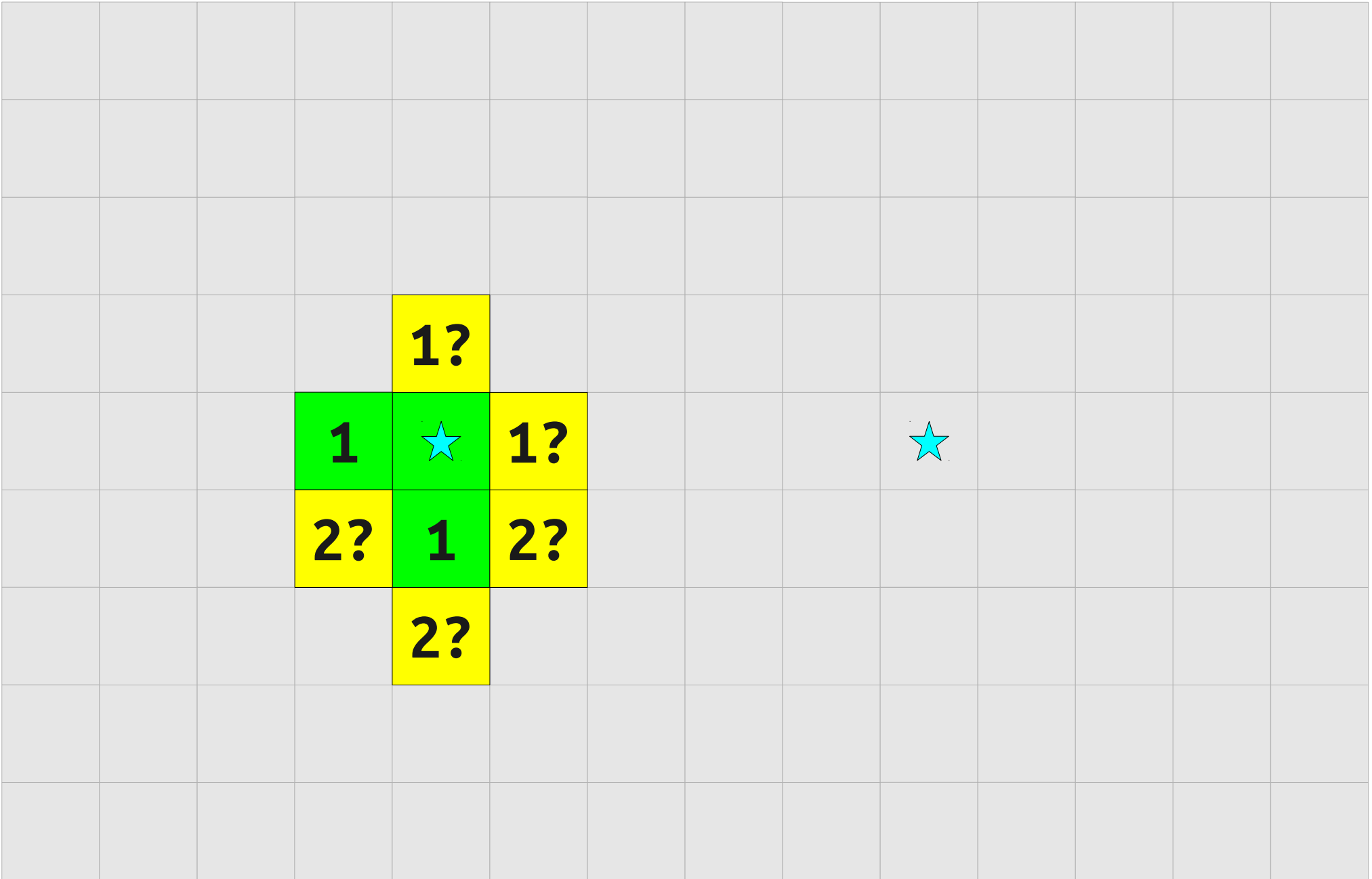


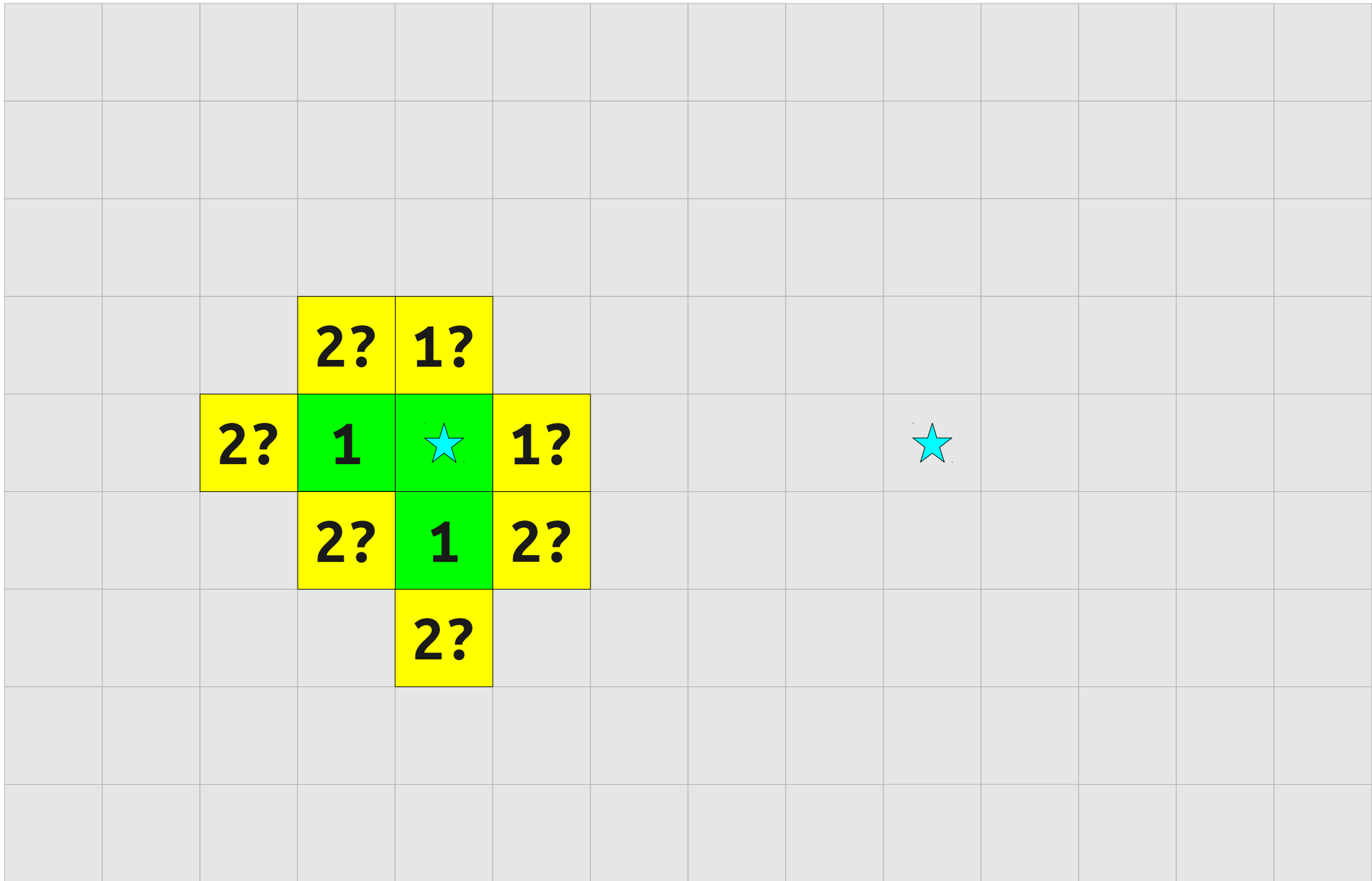


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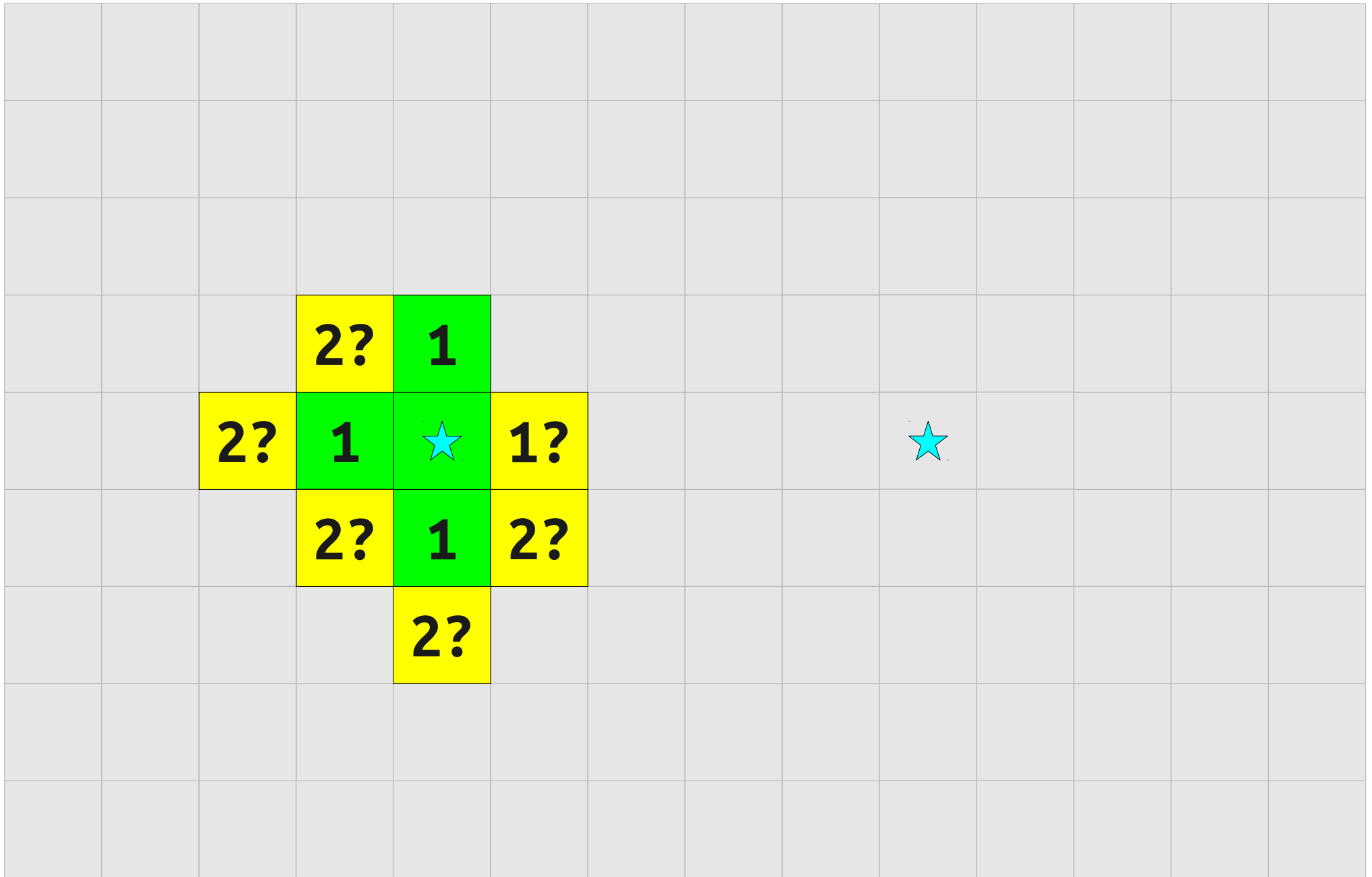


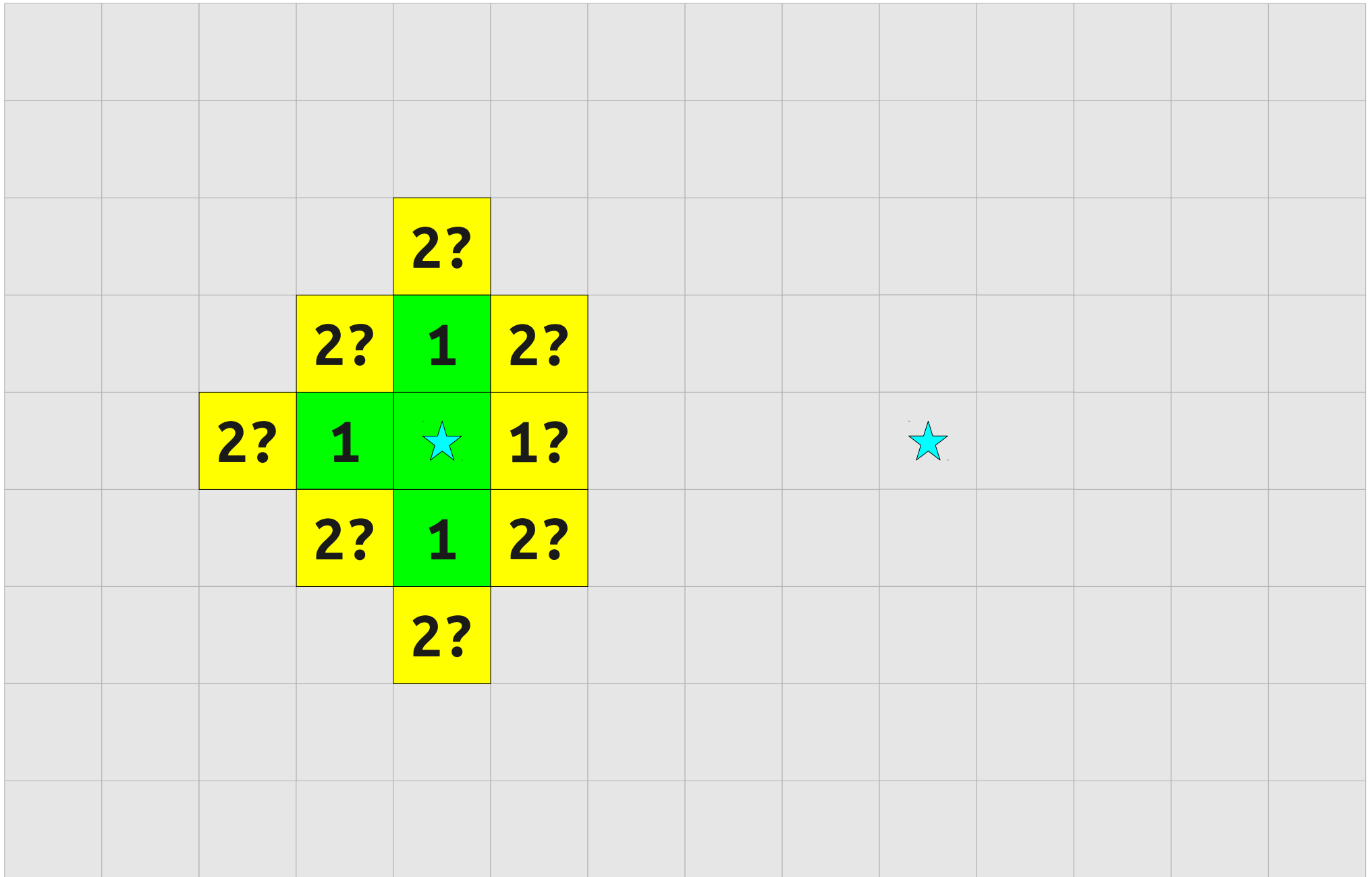


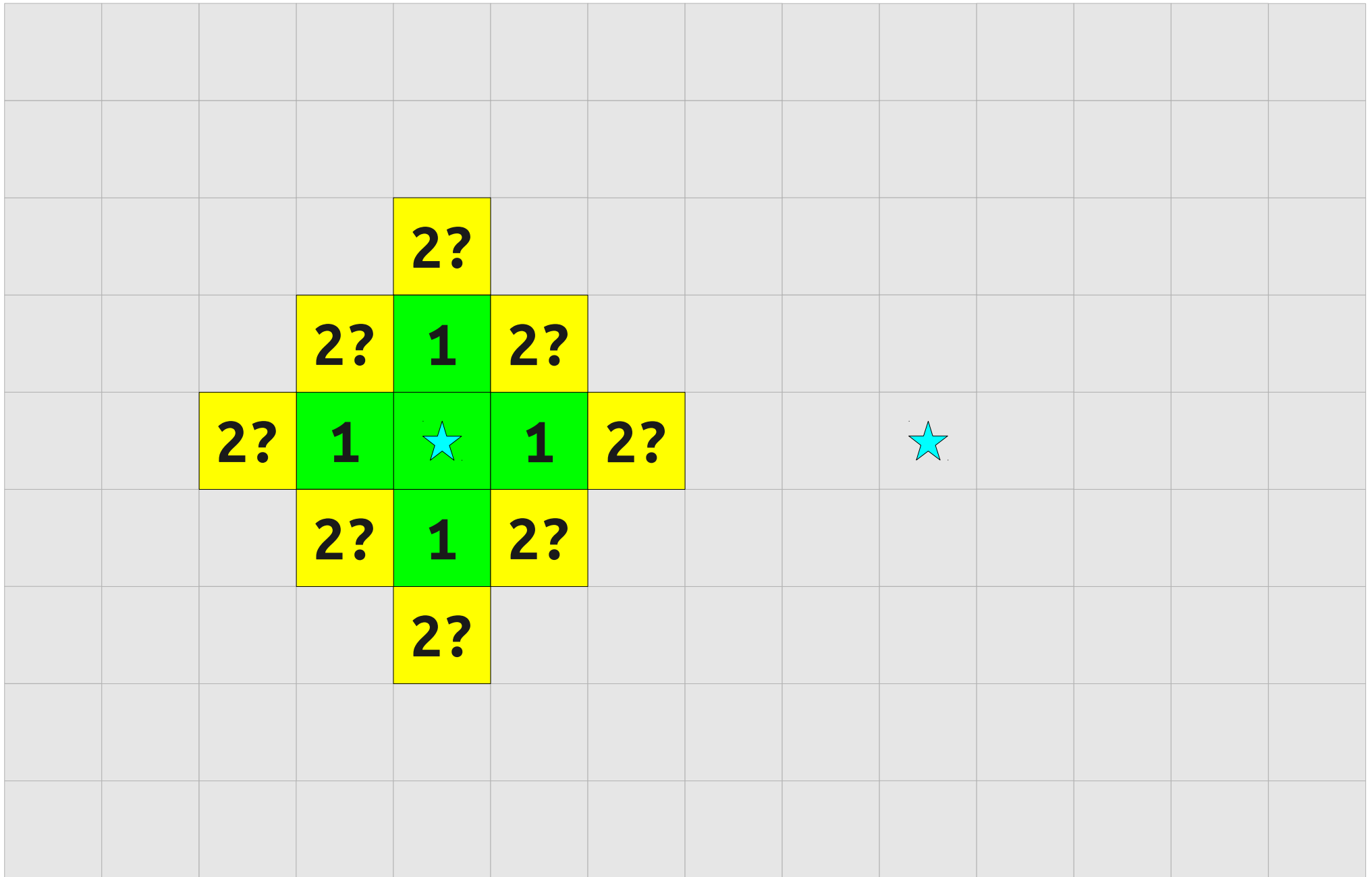


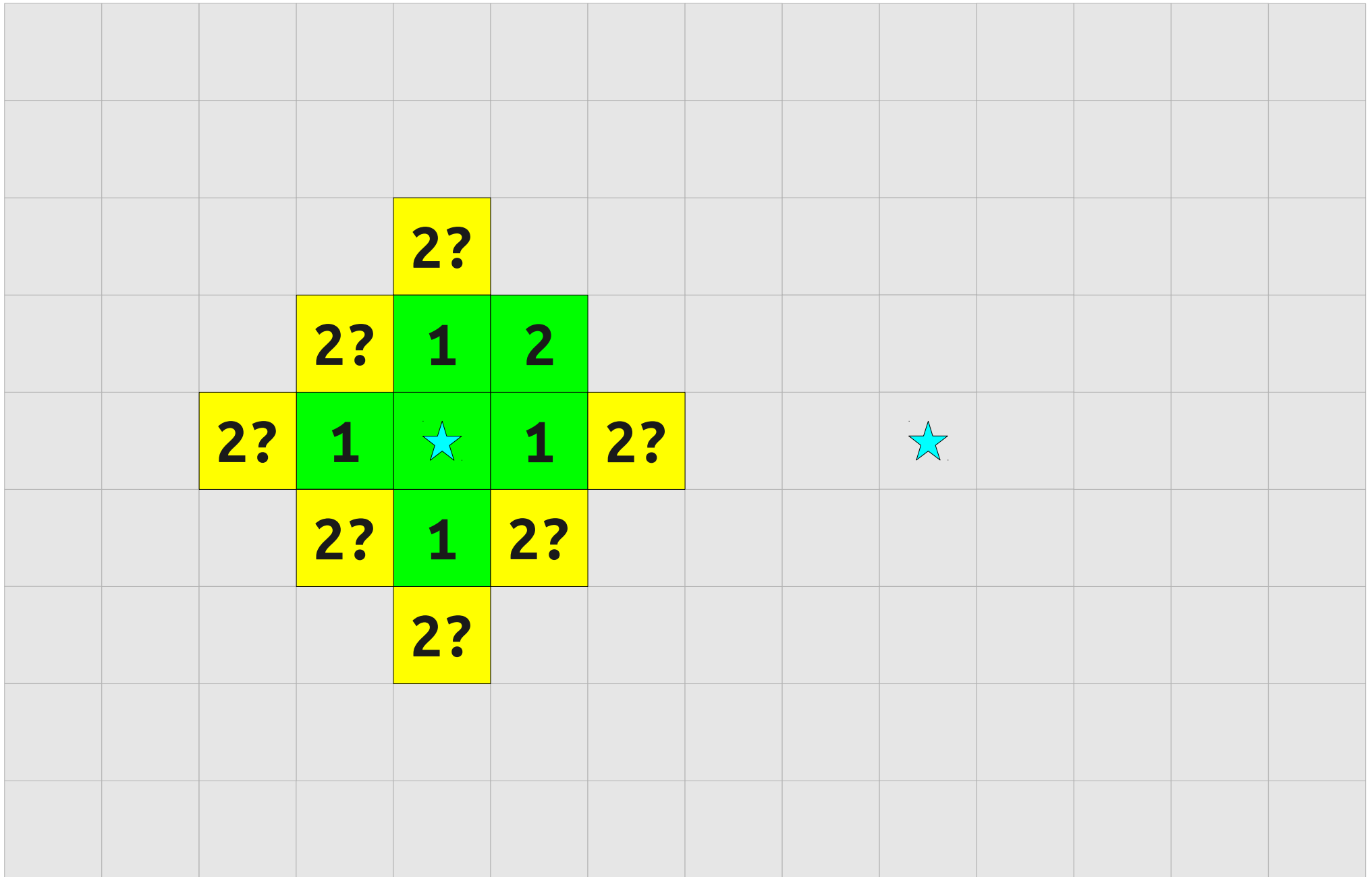


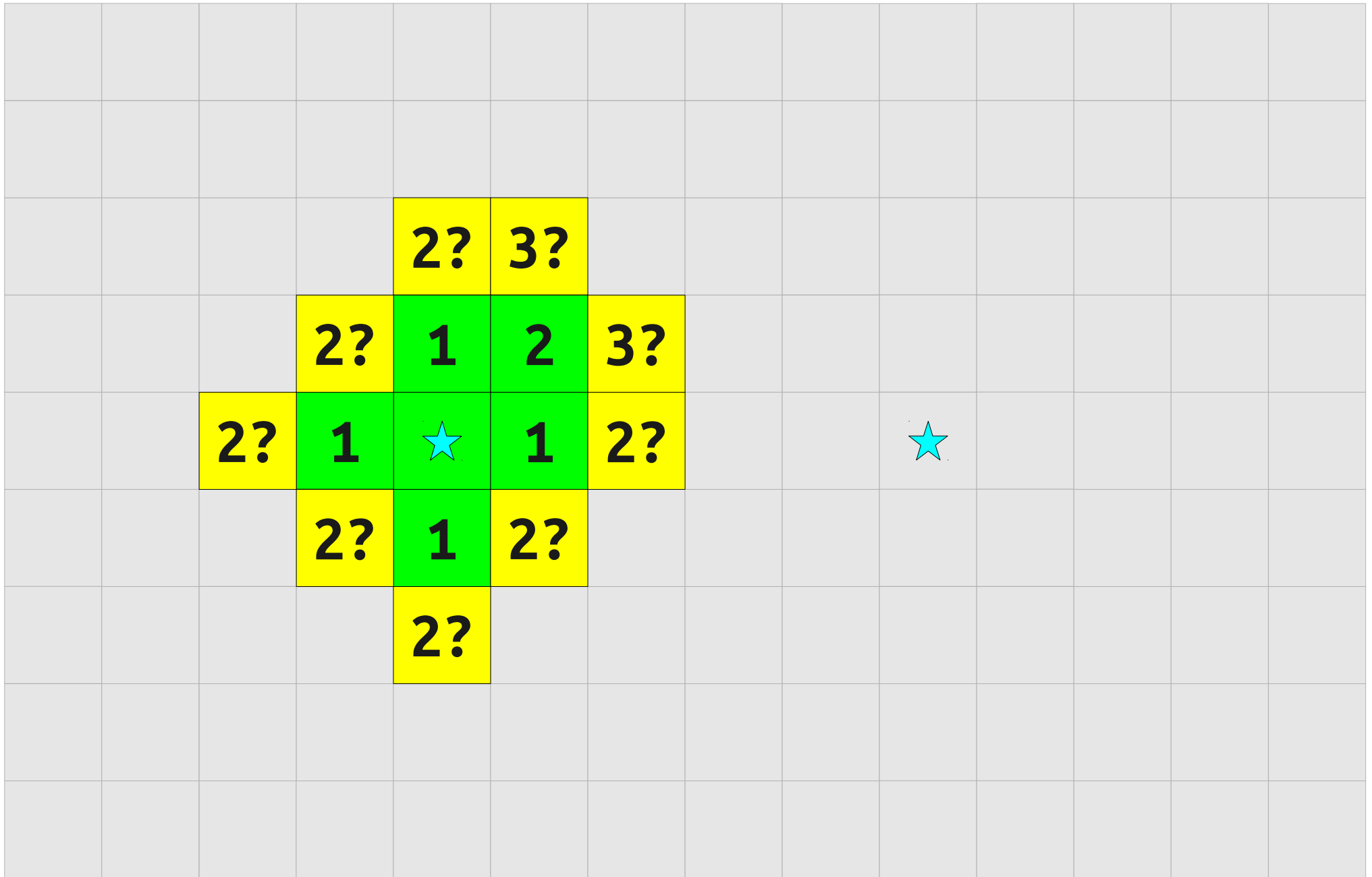


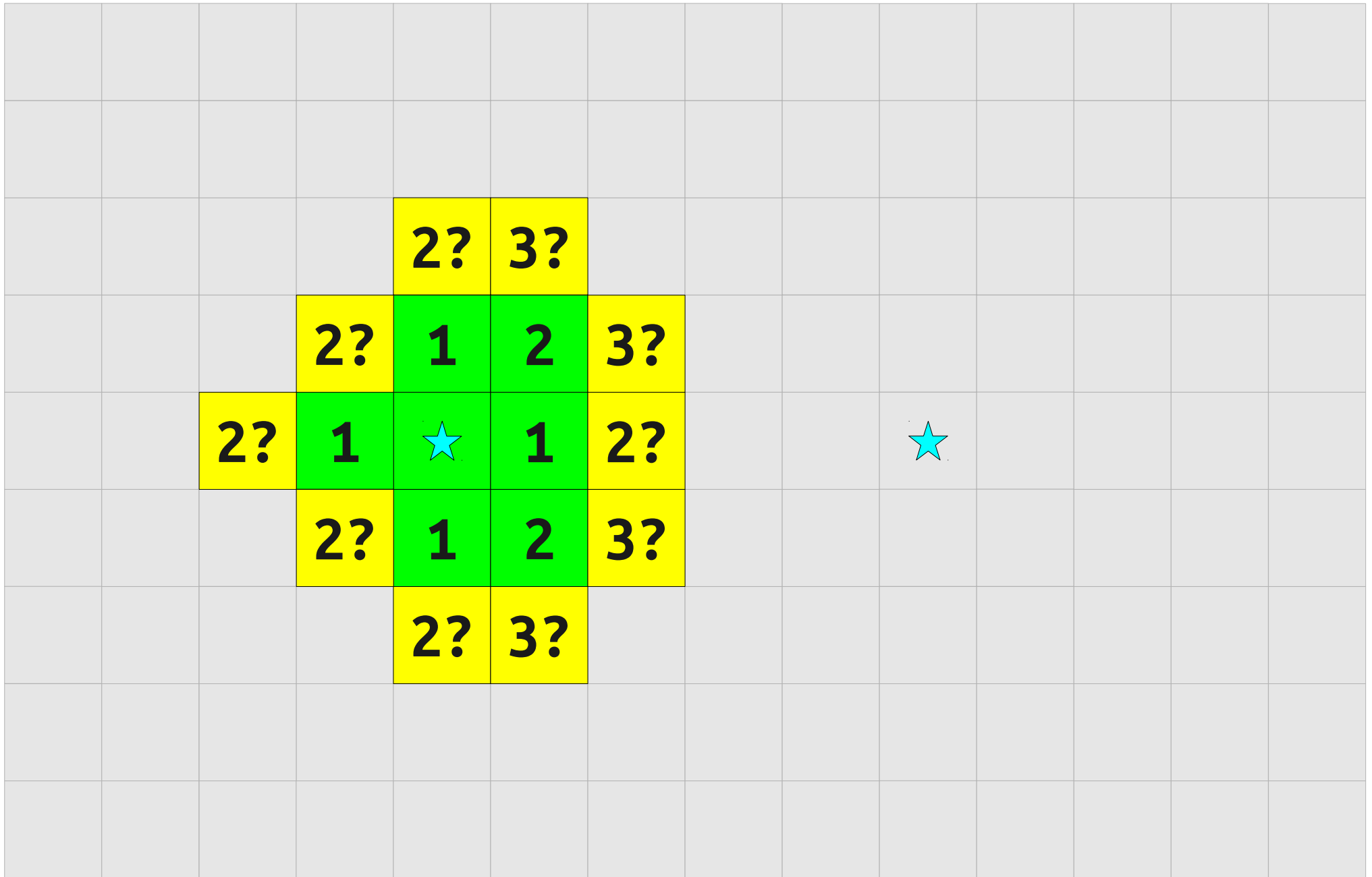


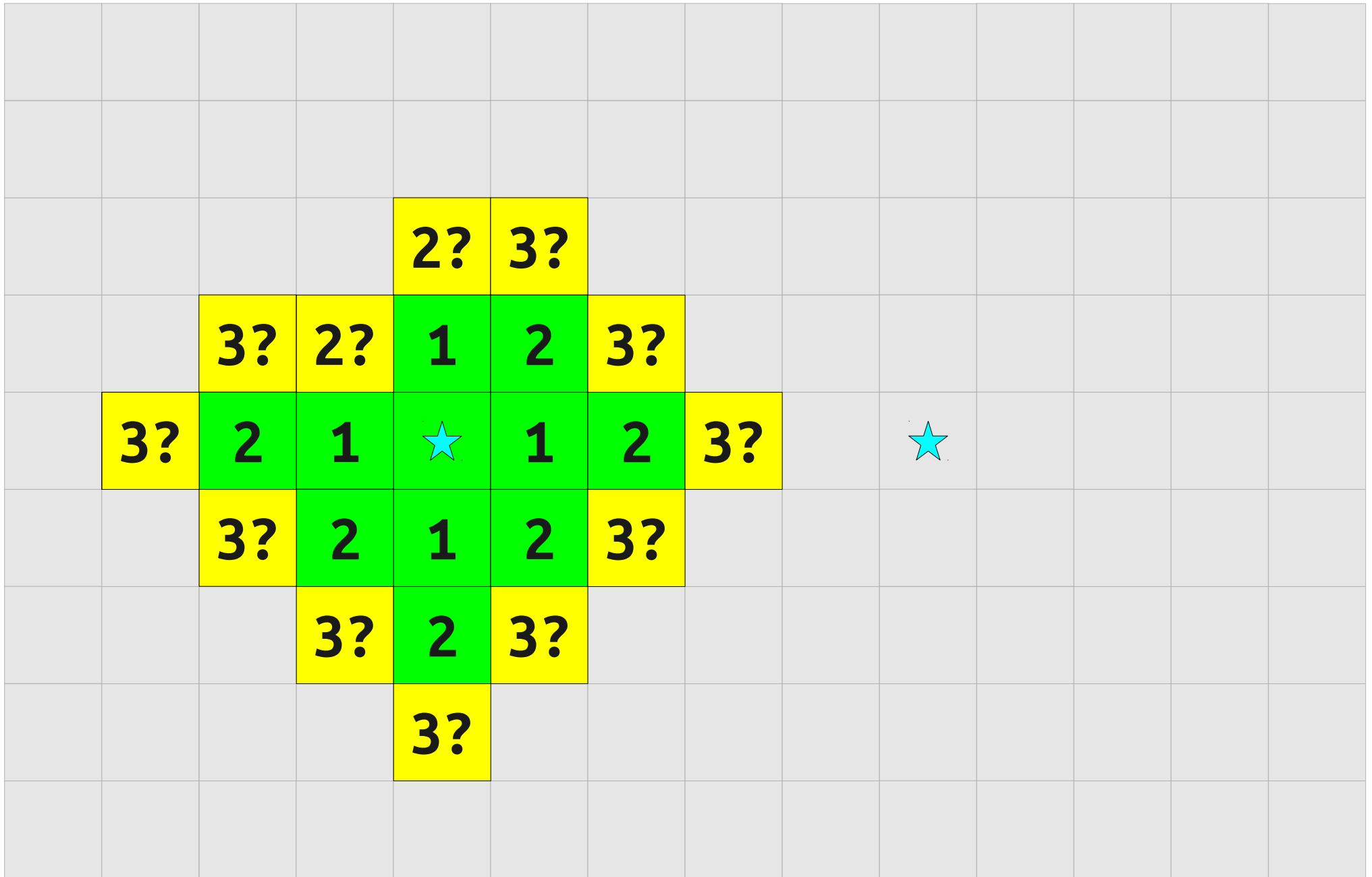


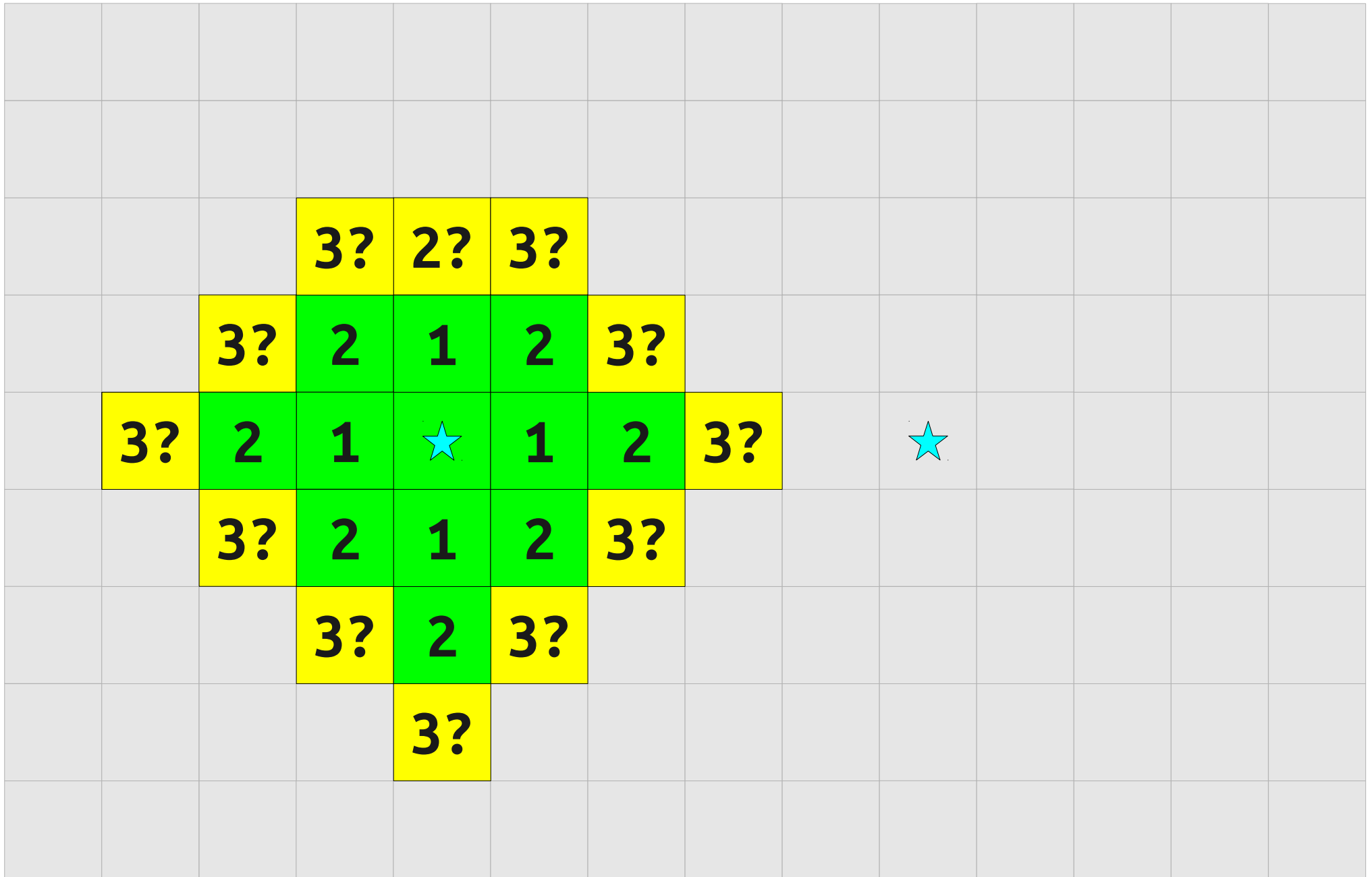


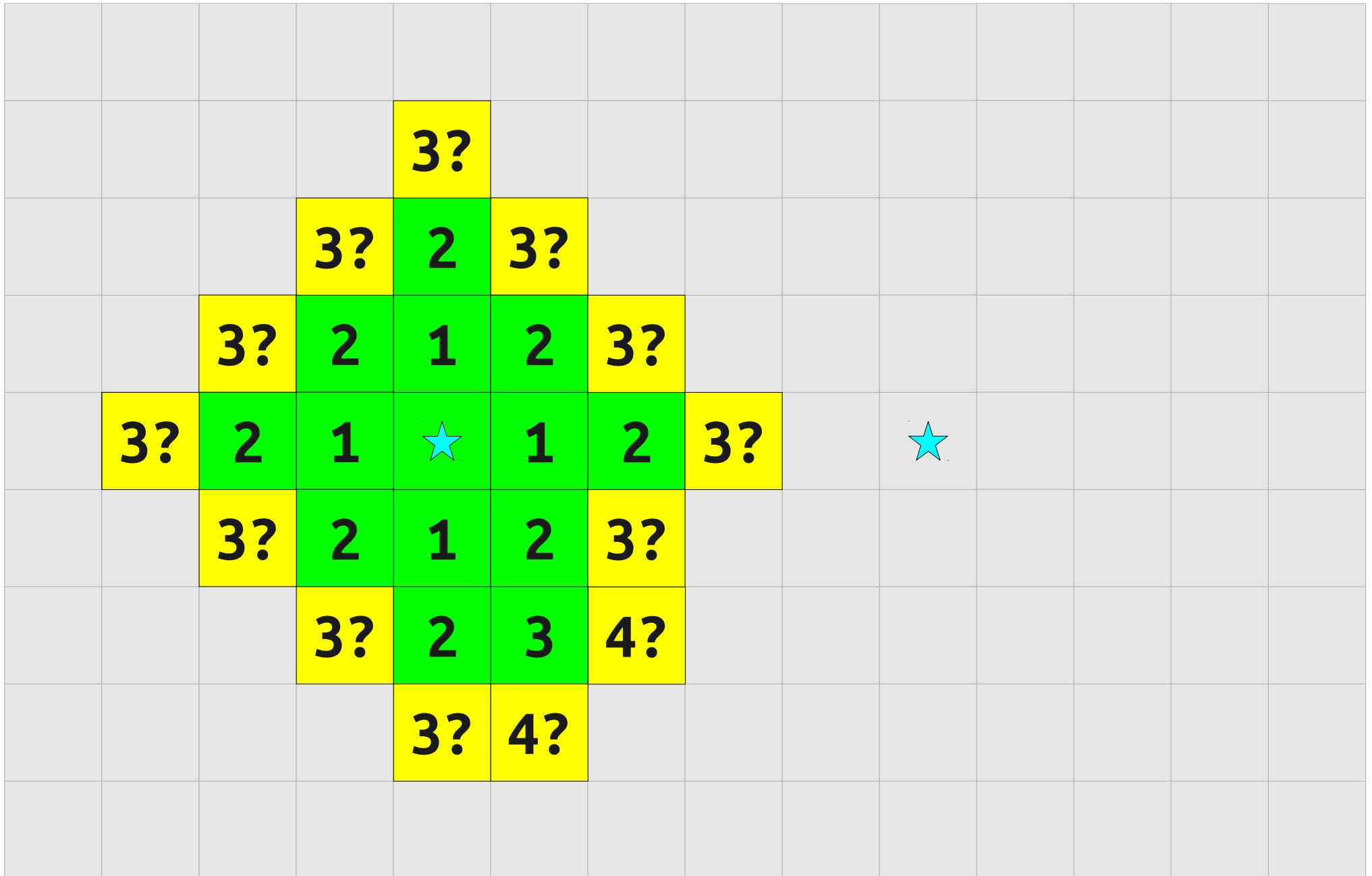


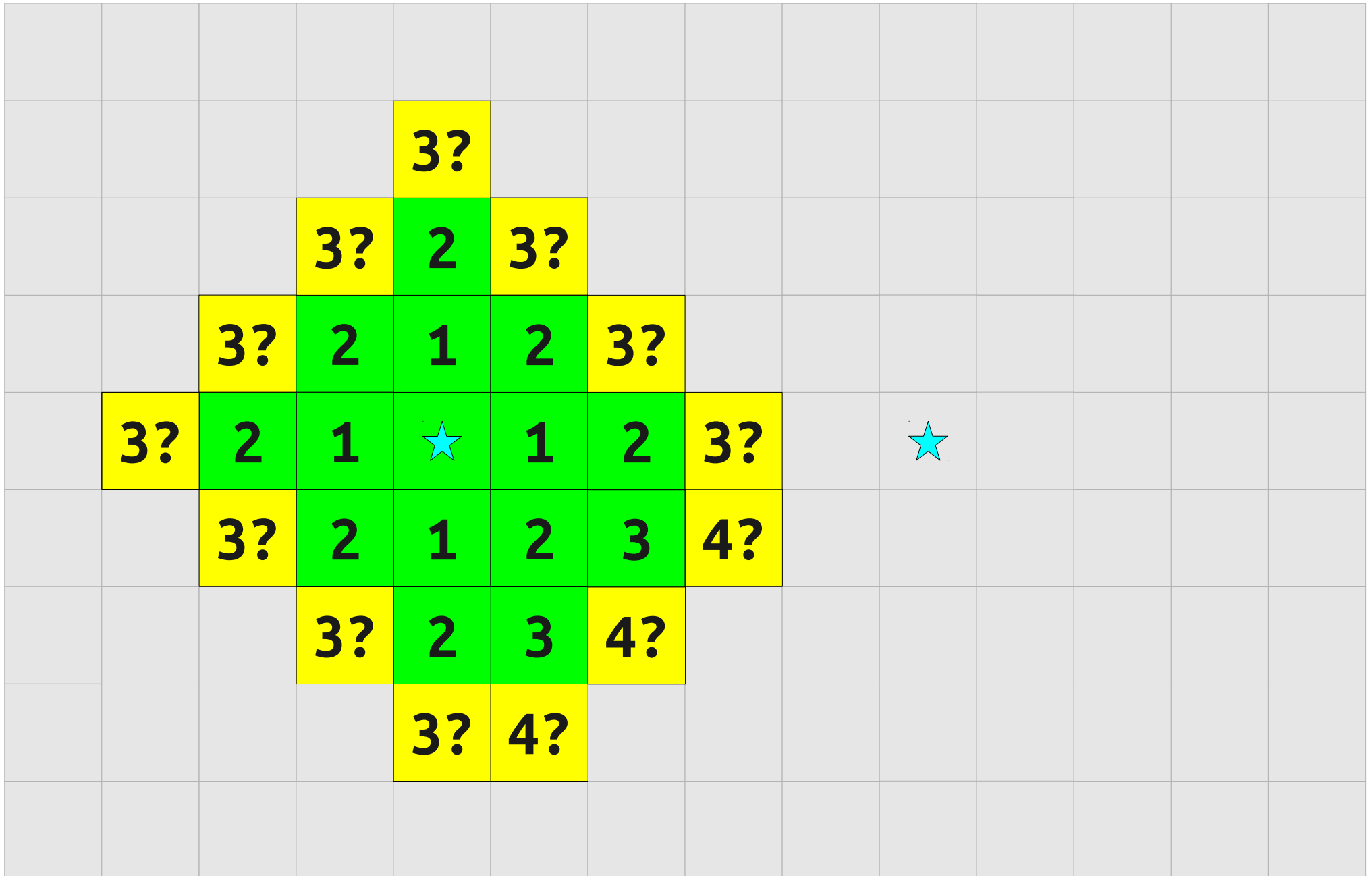


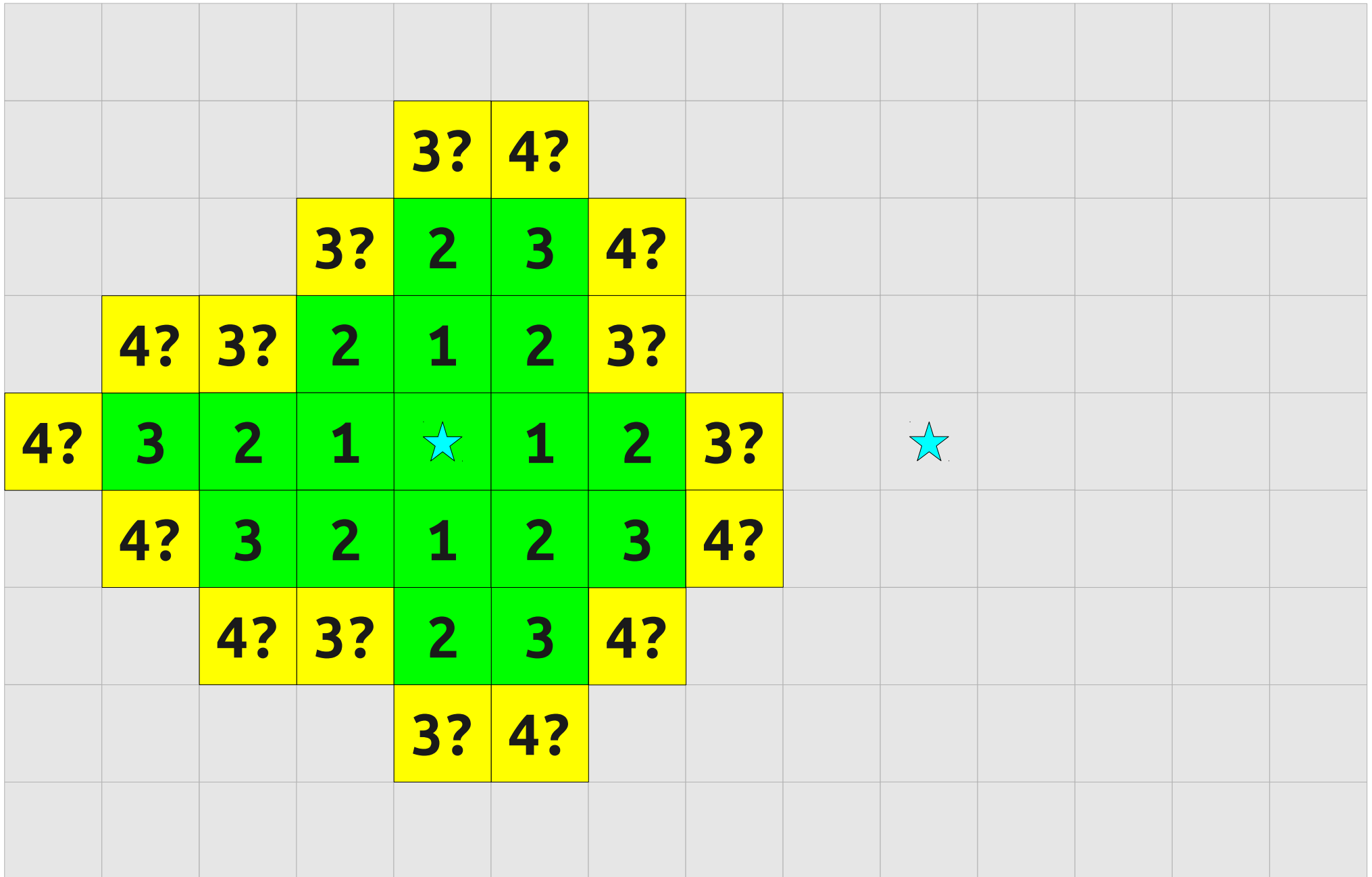


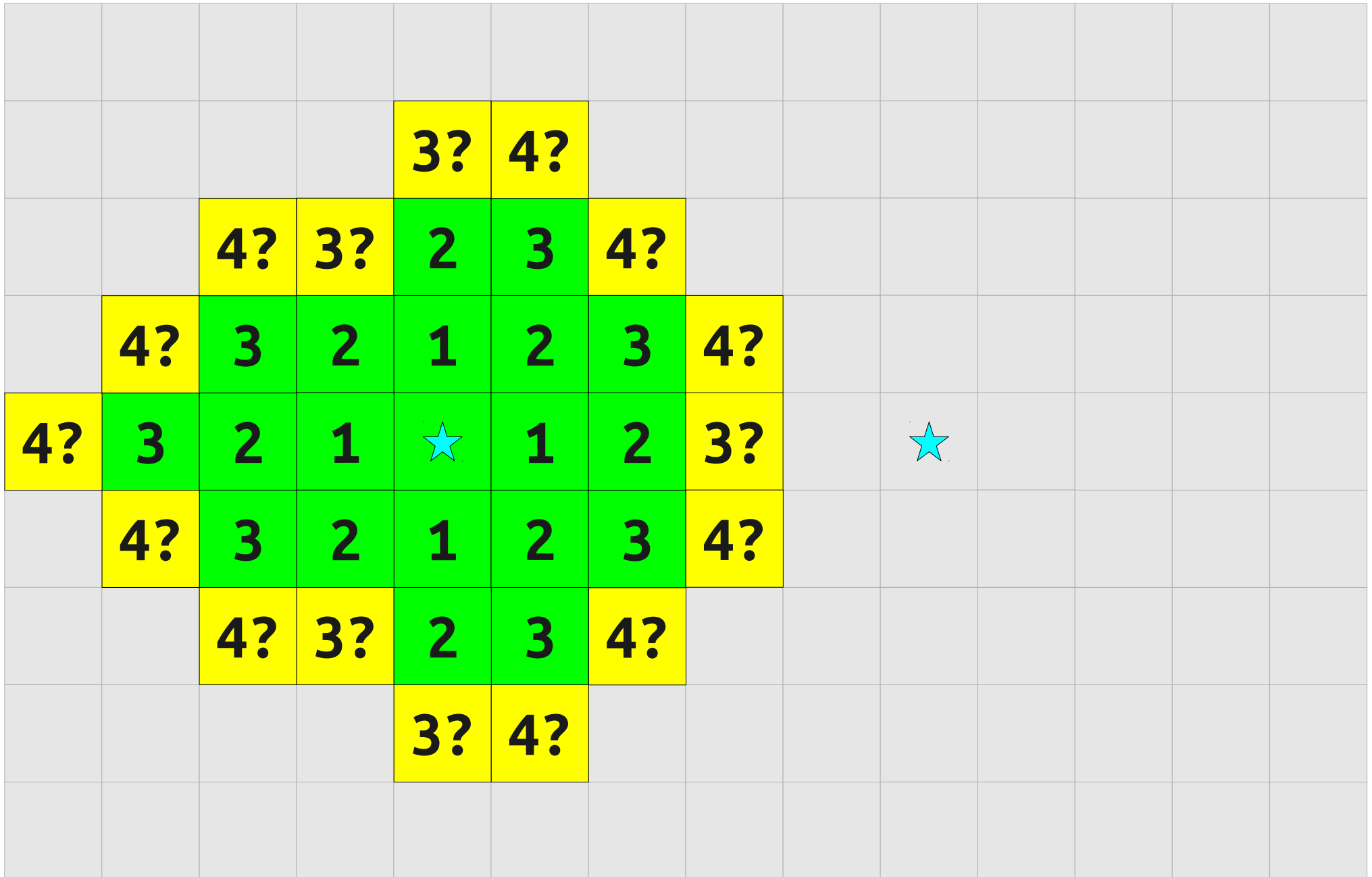


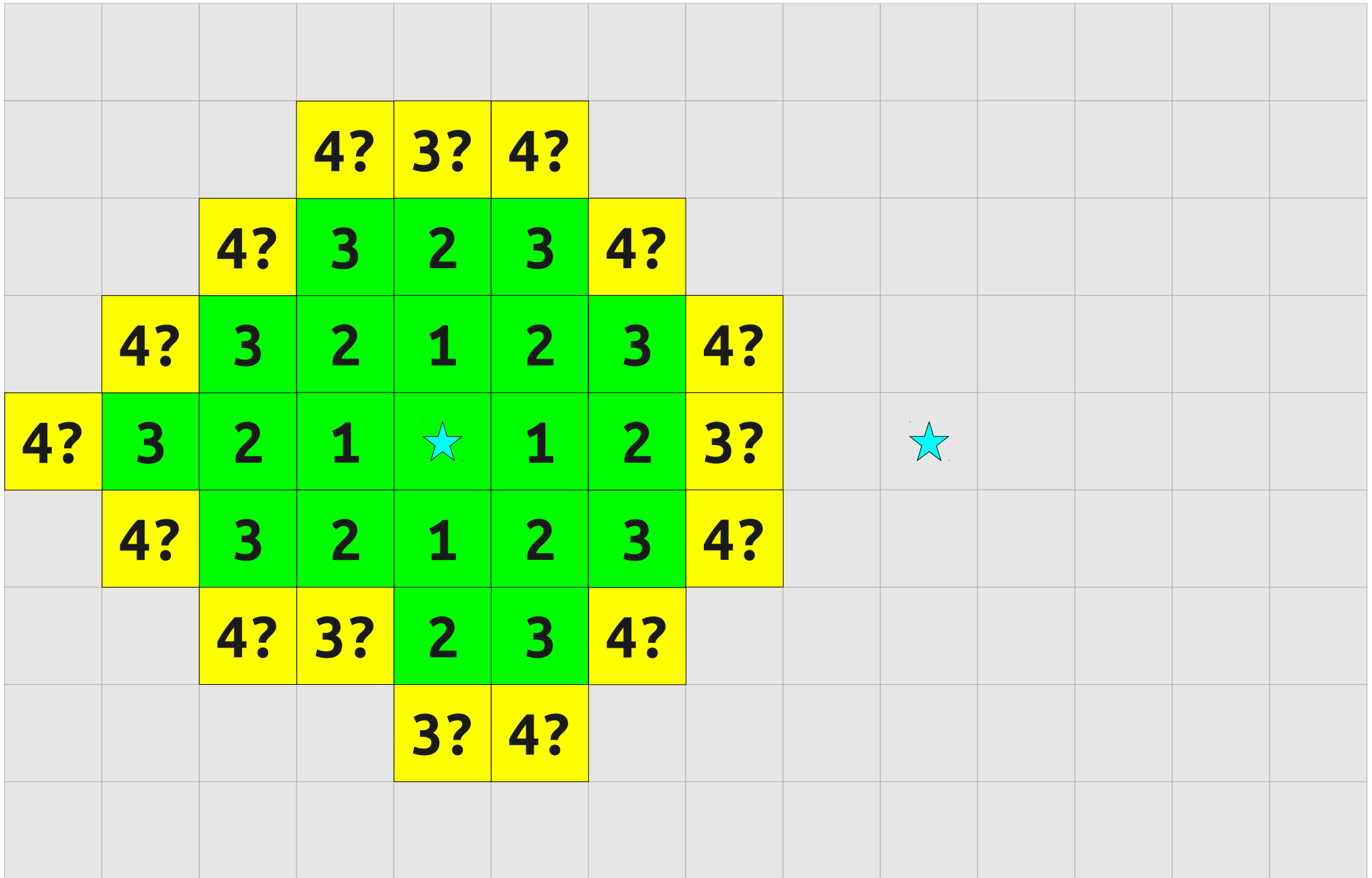


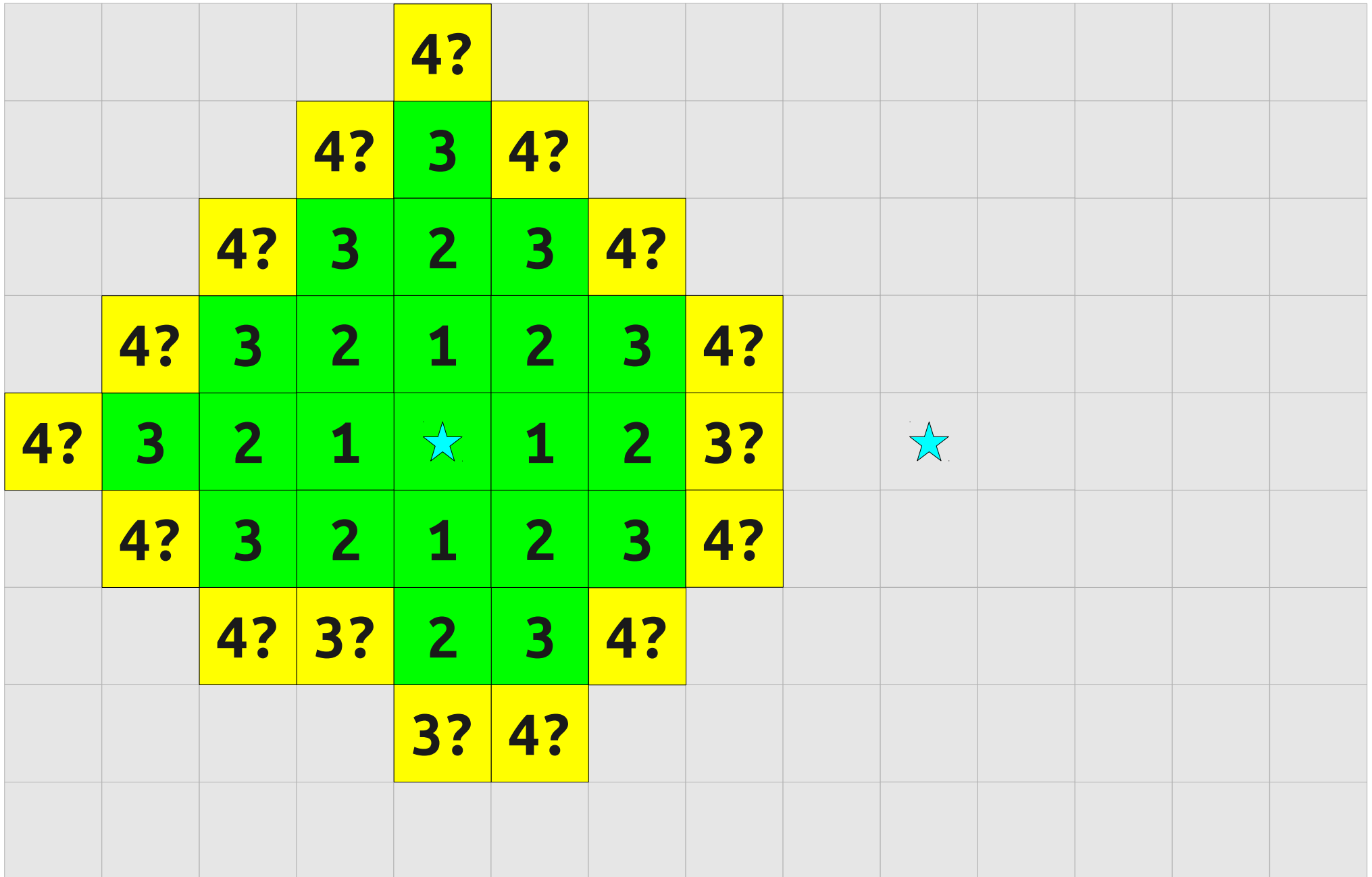








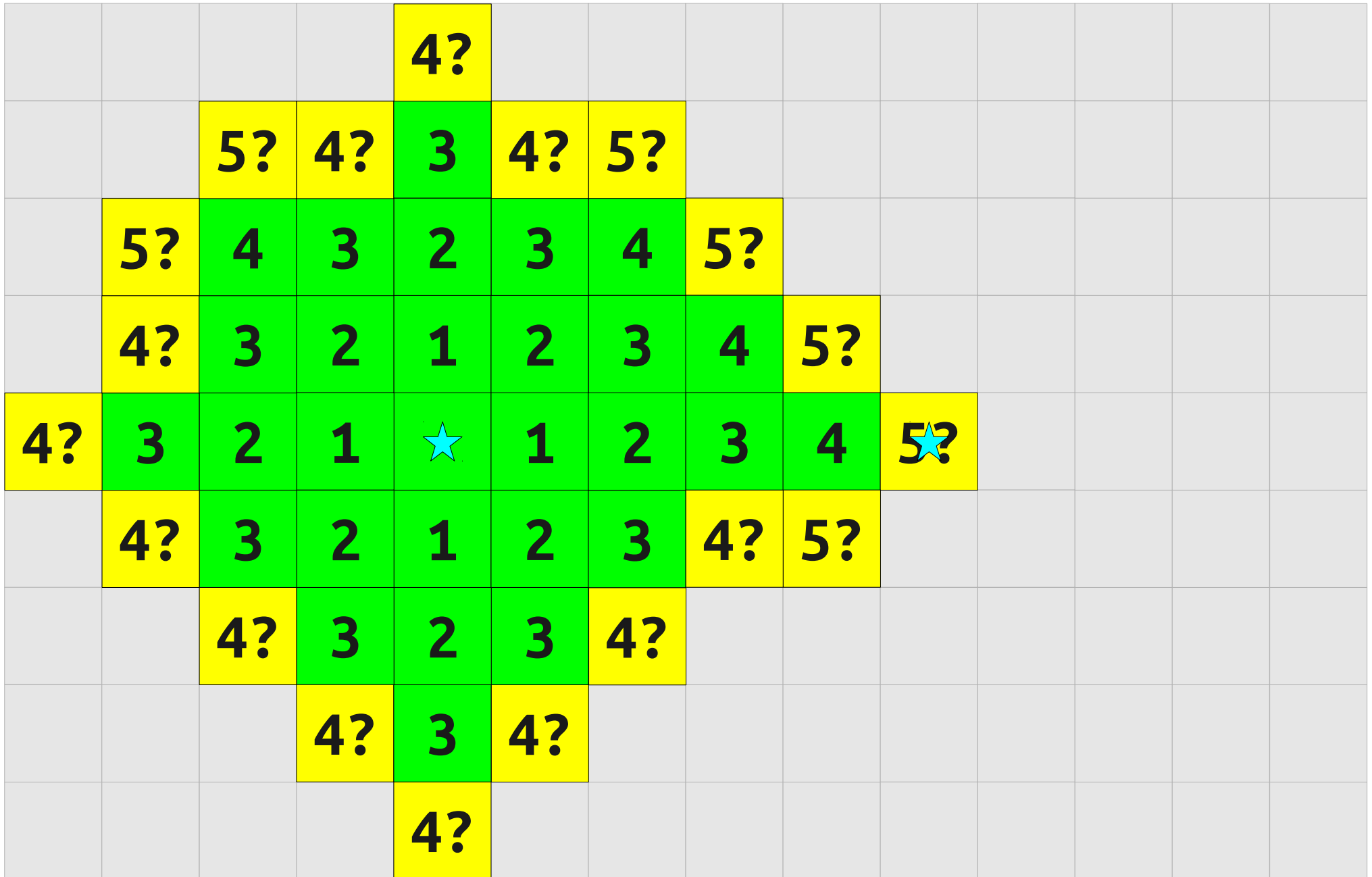


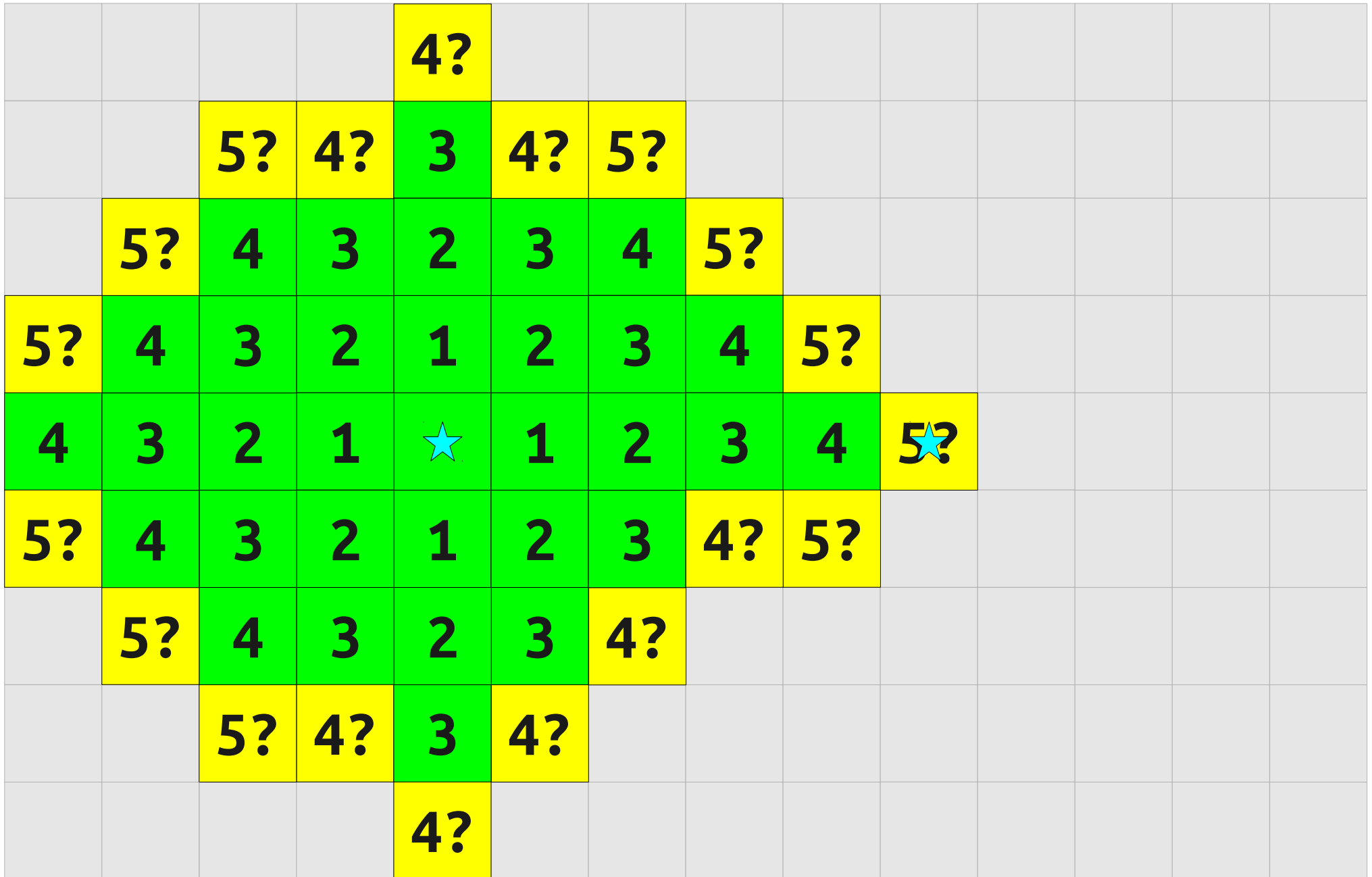


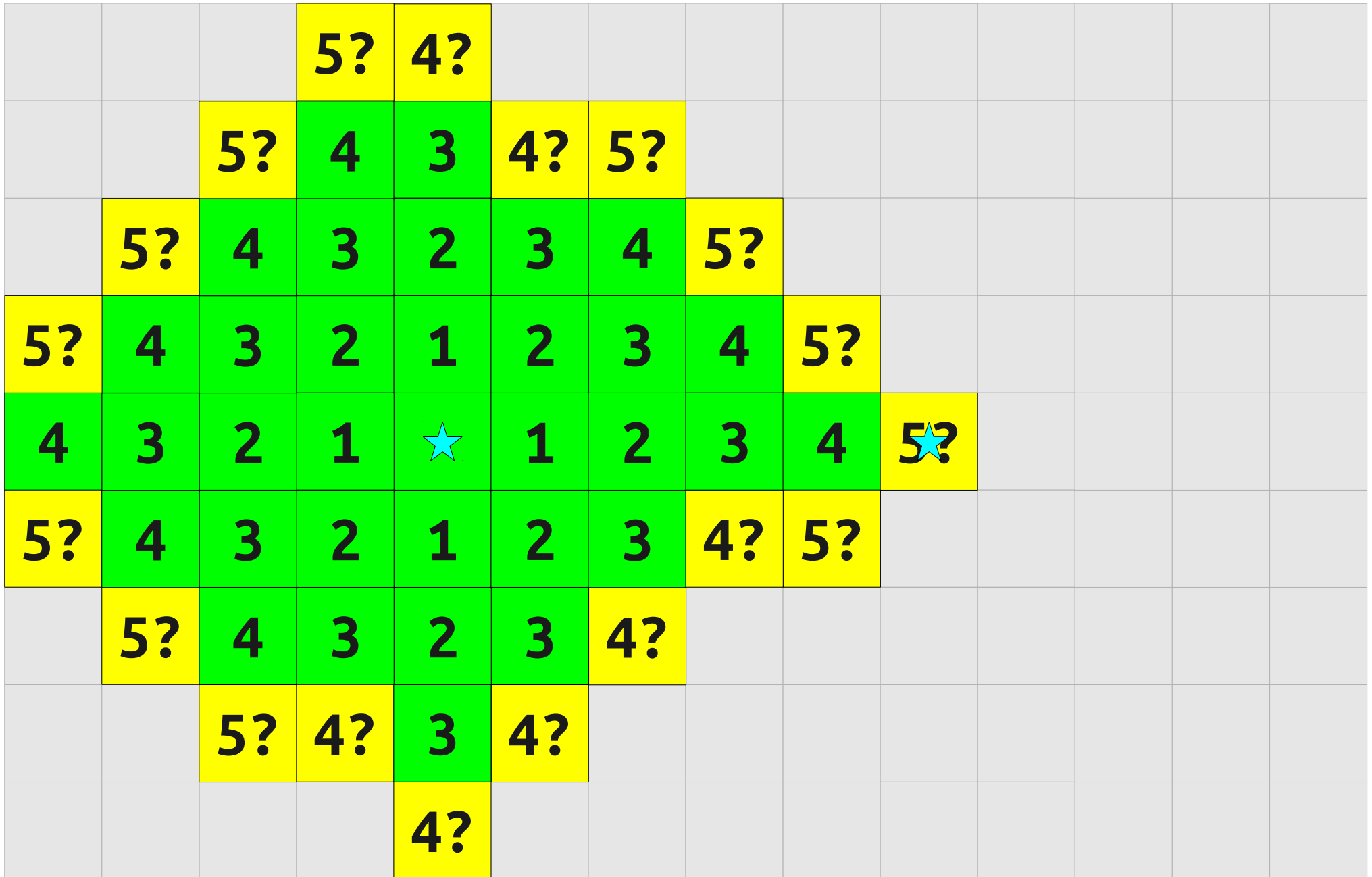




















How Dijkstra's Works

- Dijkstra's algorithm works by incrementally computing the shortest path to intermediary nodes in the graph in case they prove to be useful.
- Most of these nodes are completely in the wrong direction.
- No “big-picture” conception of how to get to the destination – the algorithm explores outward in all directions.
- Could we give the algorithm a hint?

Heuristics

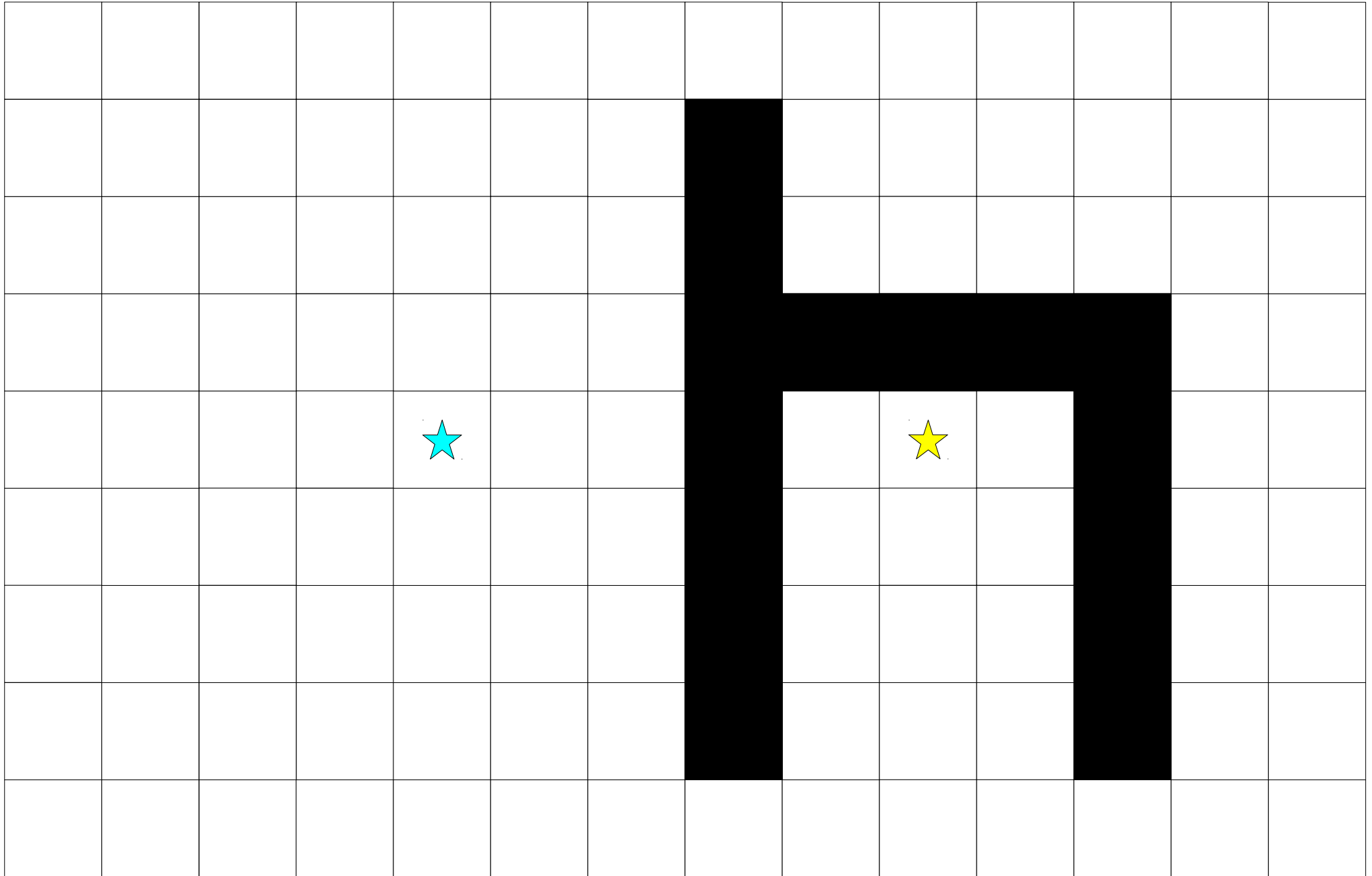
- In the context of graph searches, a **heuristic function** is a function that guesses the distance from some known node to the destination node.
- The guess doesn't have to be correct, but it should try to be as accurate as possible.
- Examples: For Google Maps, a heuristic for estimating distance might be the straight-line “as the crow flies” distance.

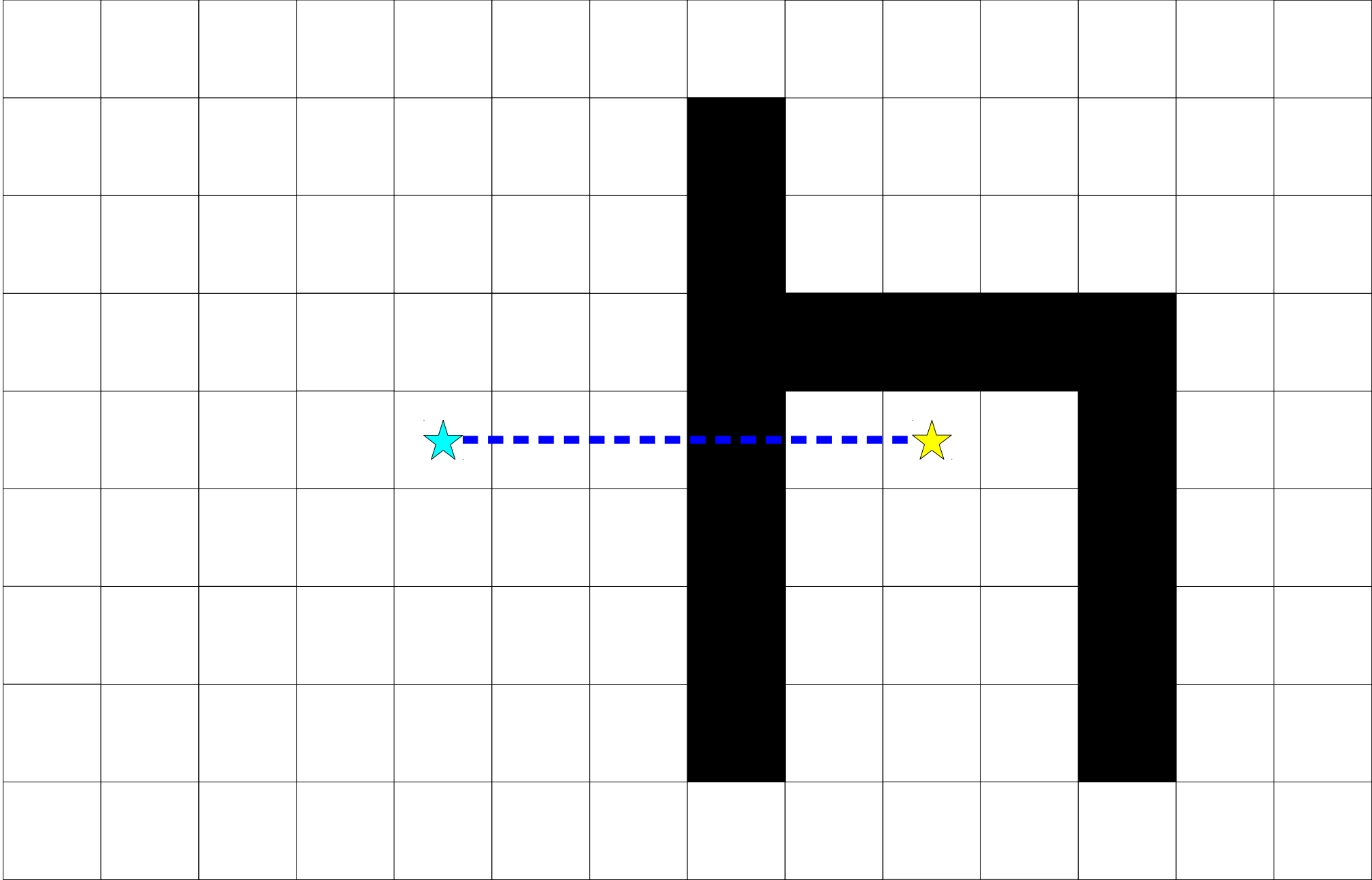
Admissible Heuristics

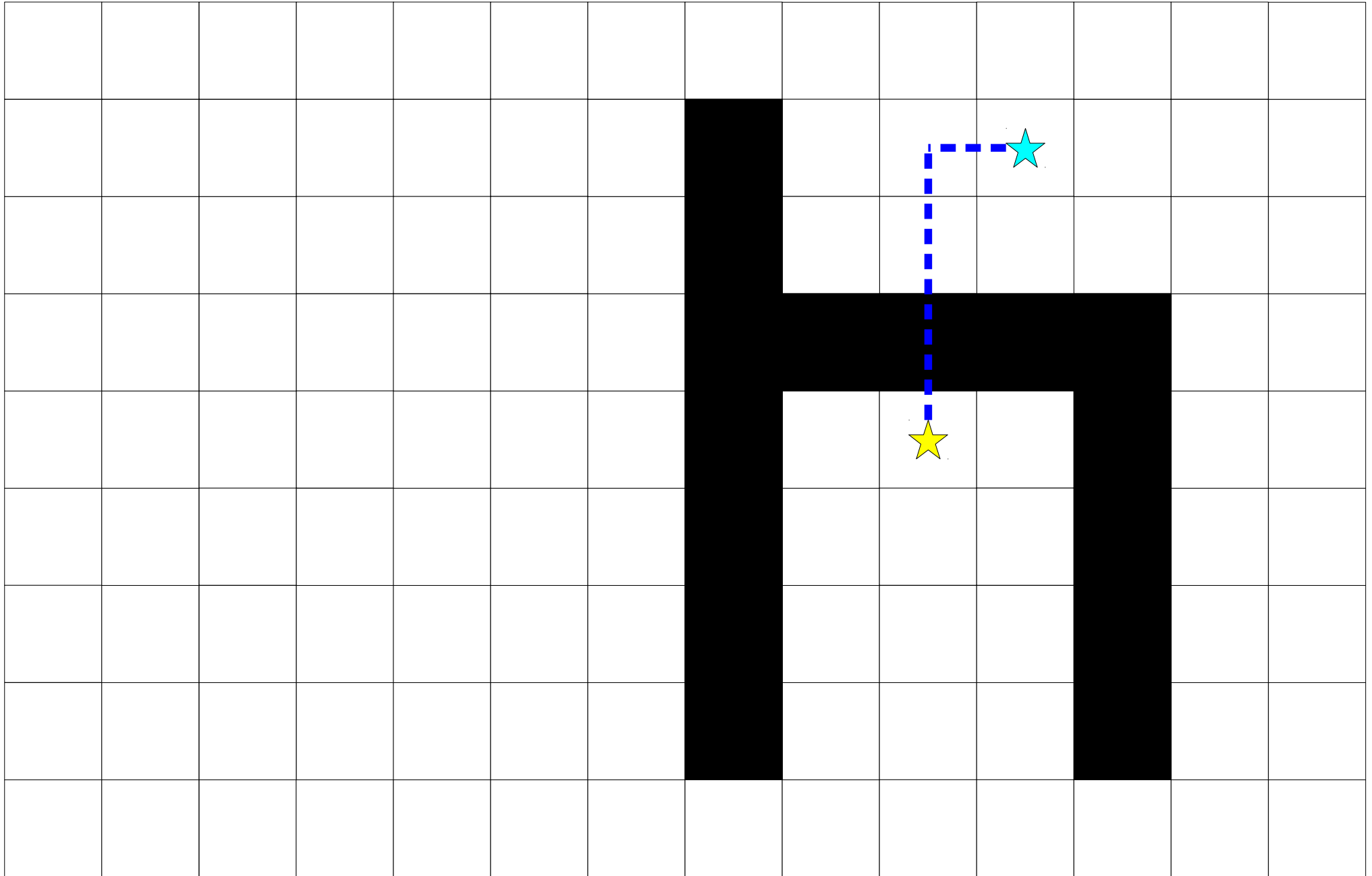
- A heuristic function is called an **admissible heuristic** if it never overestimates the distance from any node to the destination.

- In other words:

$$***predicted-distance \leq actual-distance***$$

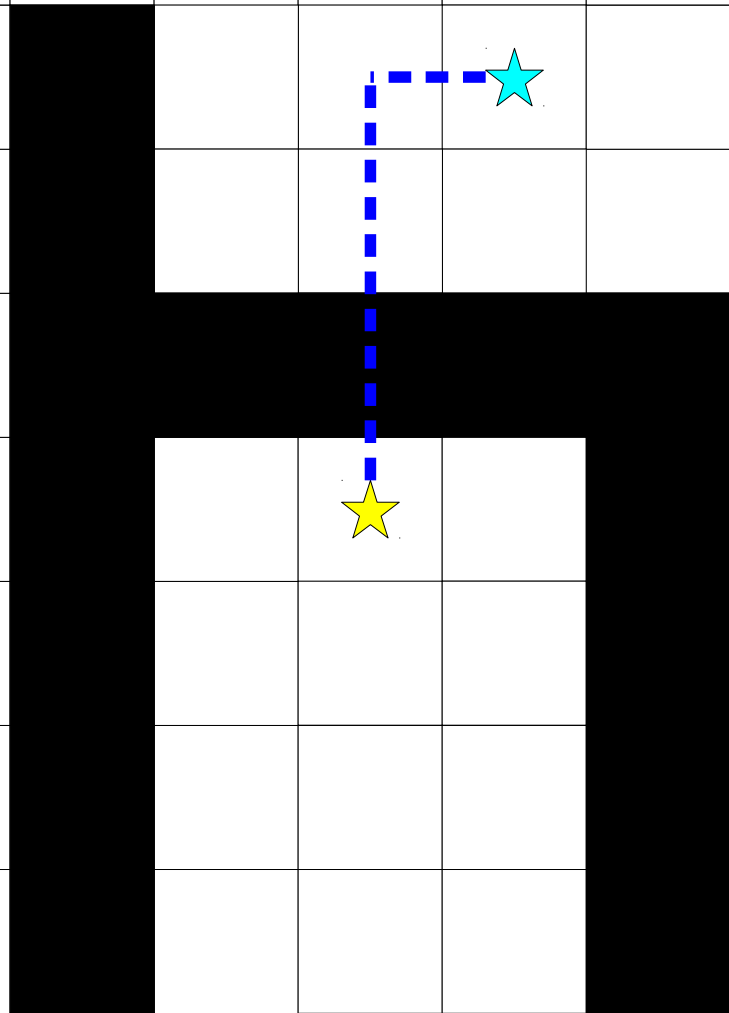






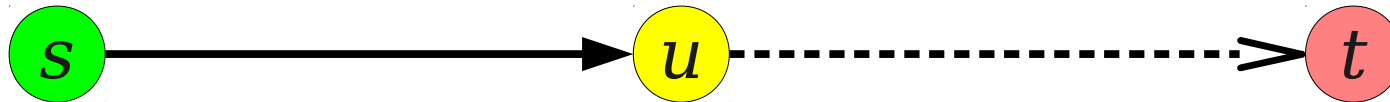
One possible heuristic:

$$|\Delta x| + |\Delta y|$$



Why Heuristics Matter

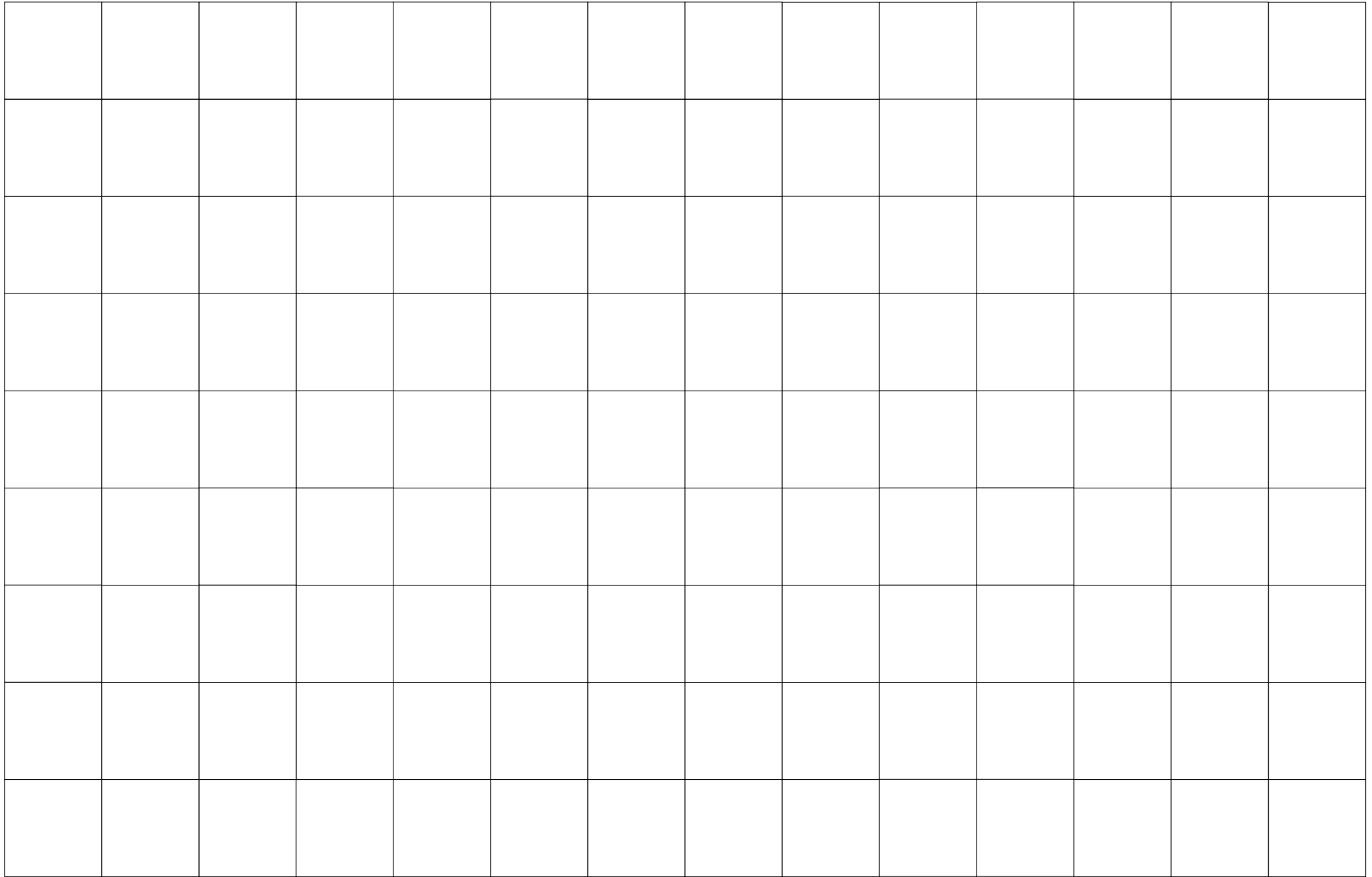
- We can modify Dijkstra's algorithm by introducing heuristic functions.
- Given any node u , there are two associated costs:

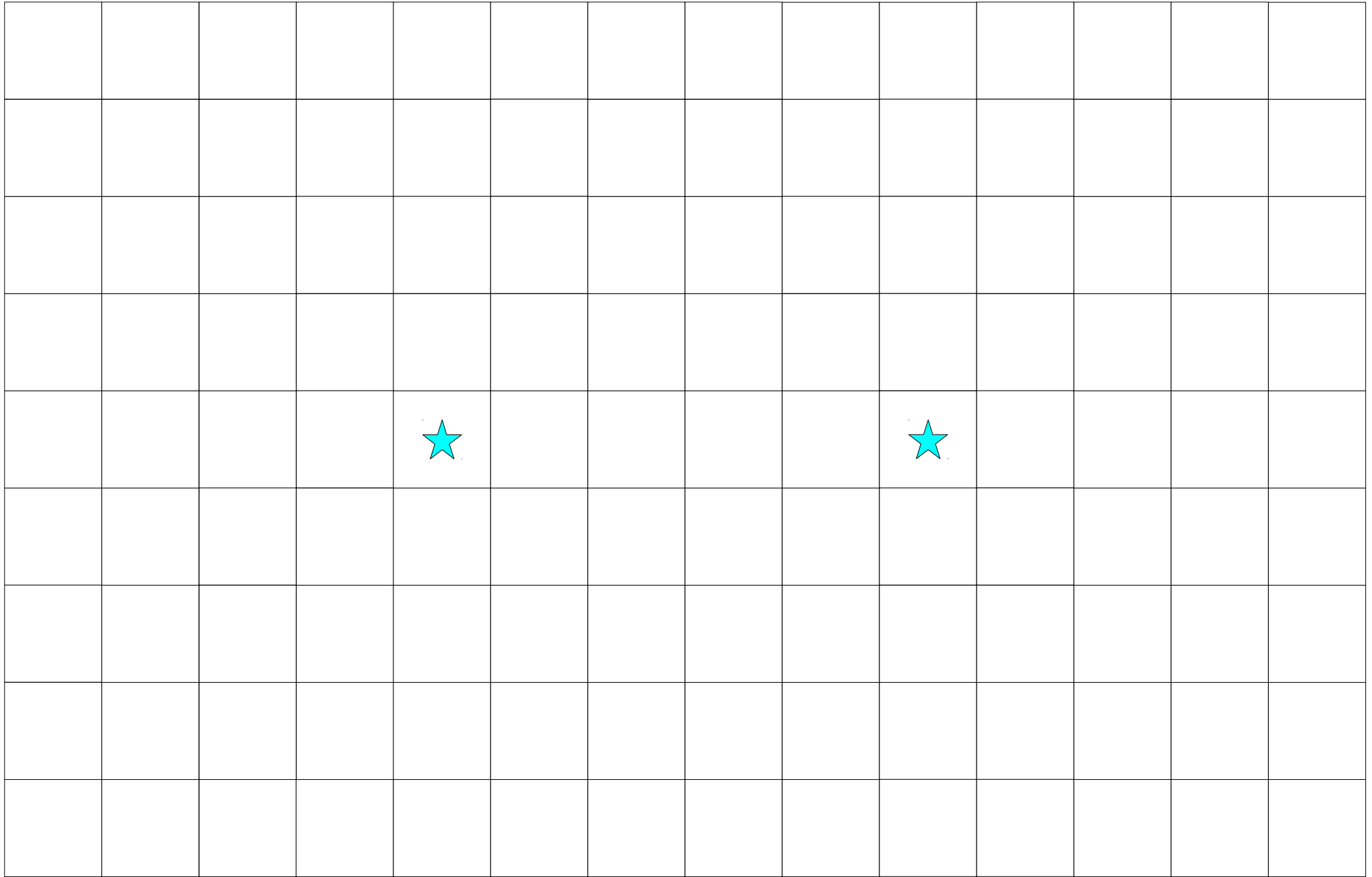


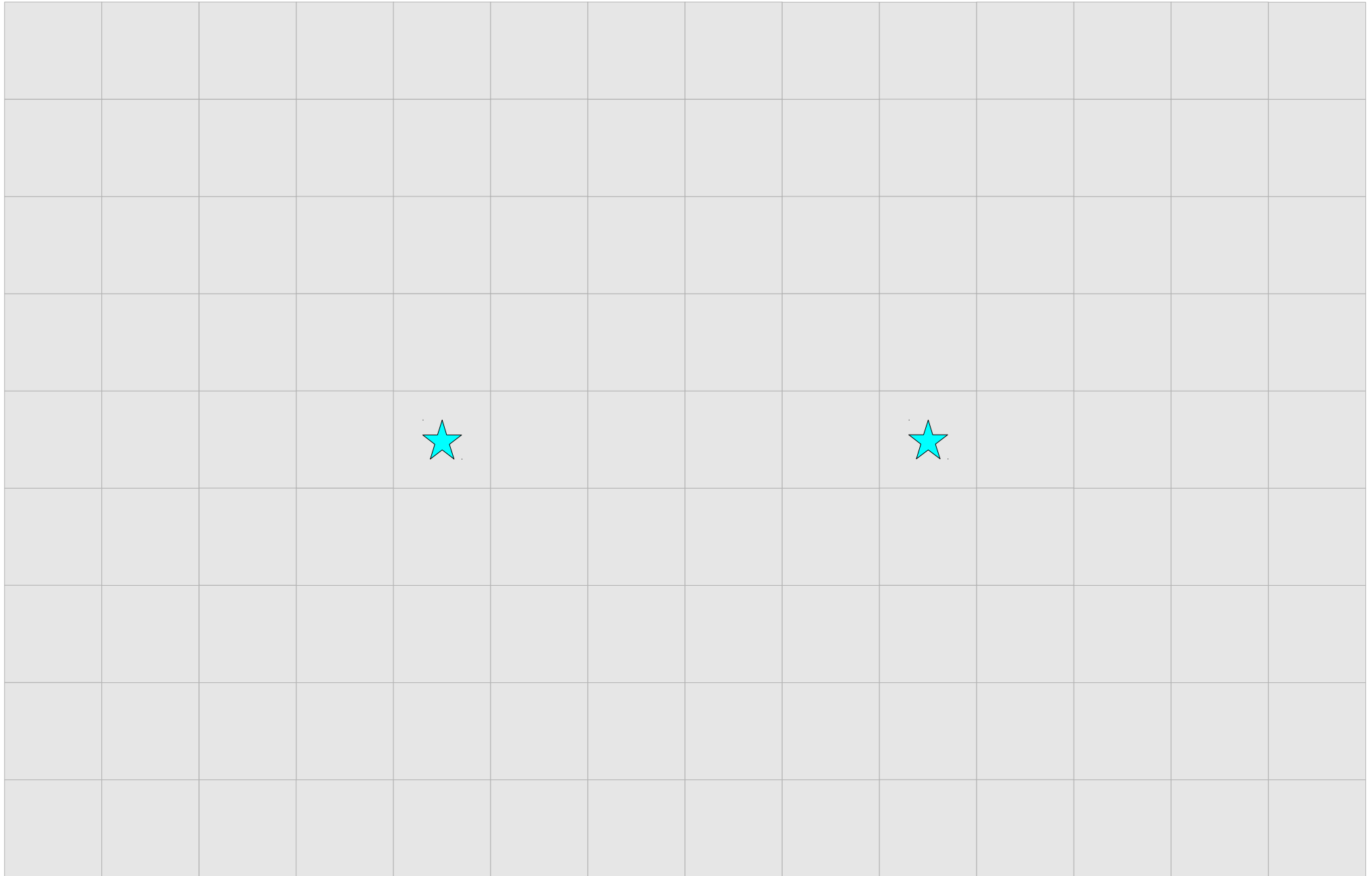
- The actual distance from the start node s .
- The heuristic distance from u to the end node t .
- Key idea: Run Dijkstra's algorithm, but use the following priority in the priority queue:

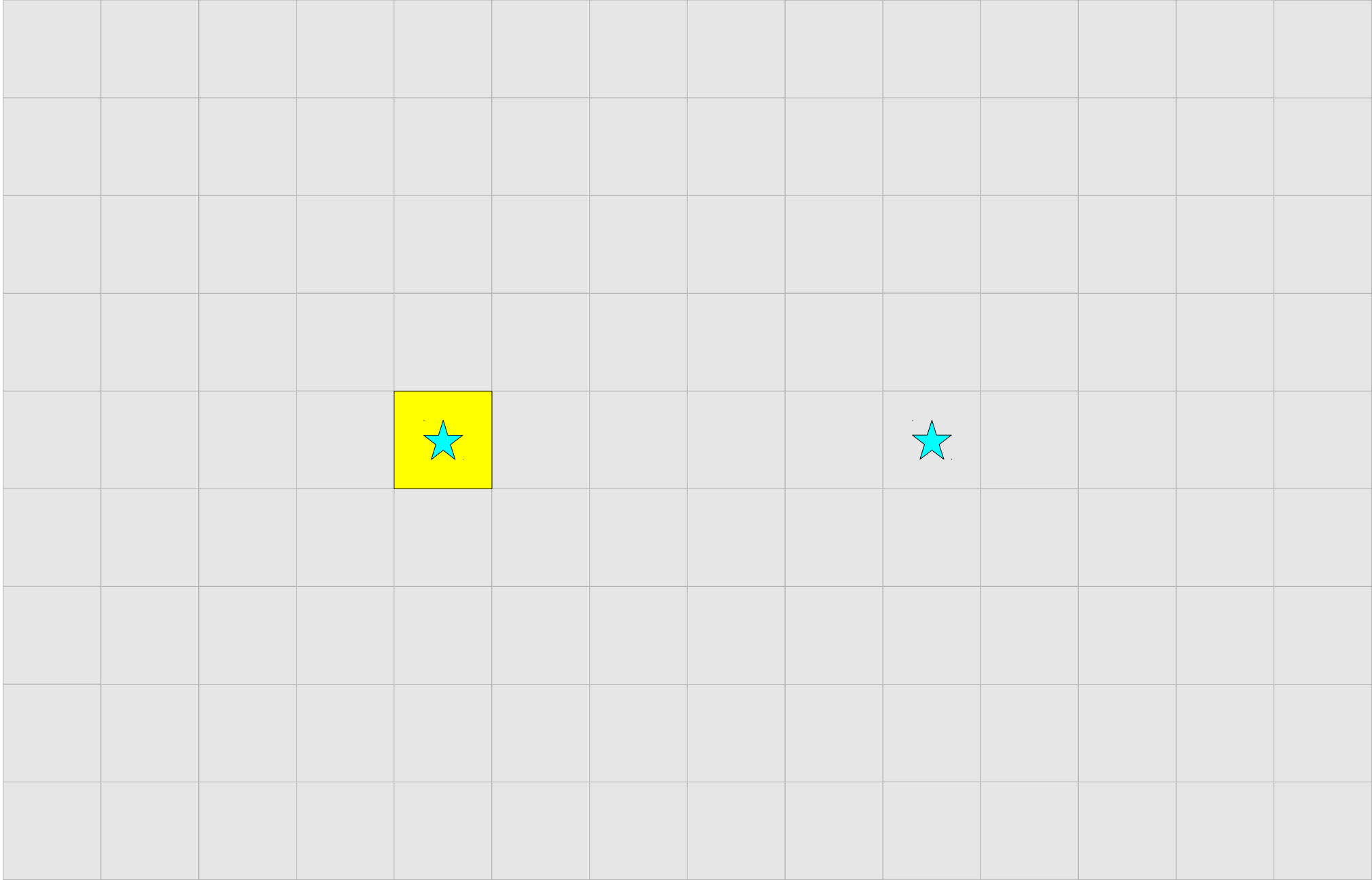
$$\mathbf{priority}(u) = \mathbf{distance}(s, u) + \mathbf{heuristic}(u, t)$$

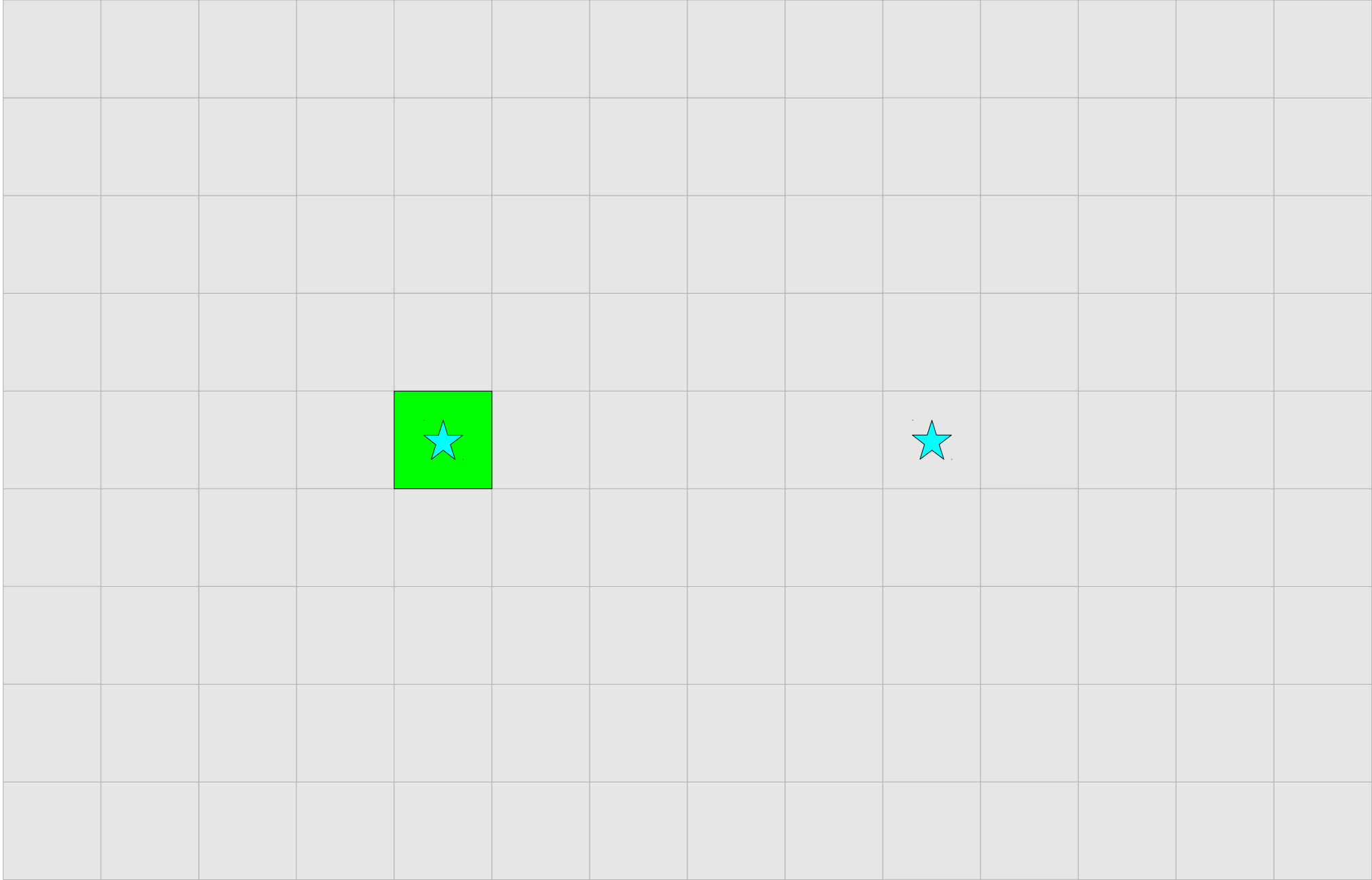
- This modification of Dijkstra's algorithm is called the **A* search algorithm**.

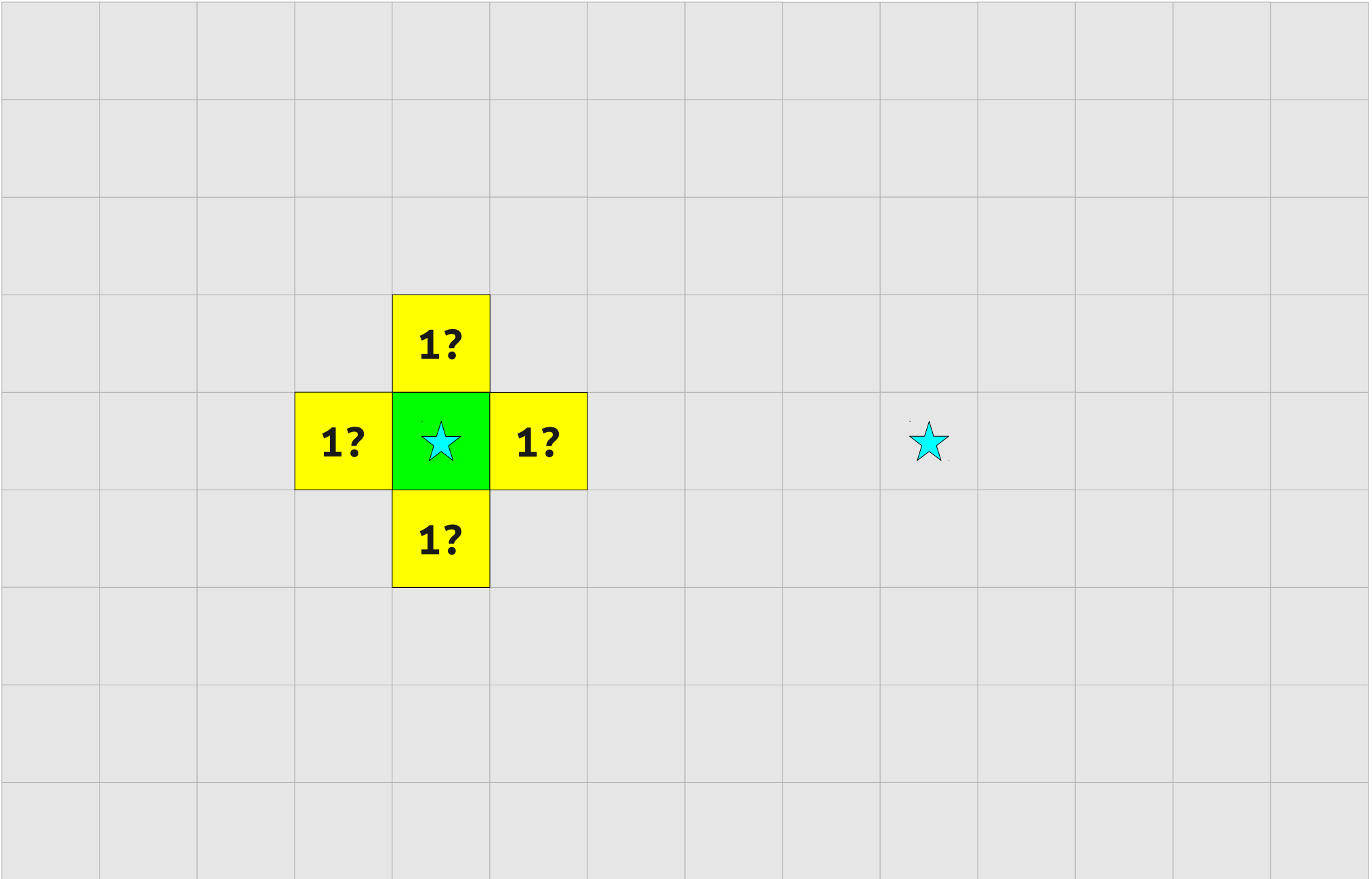


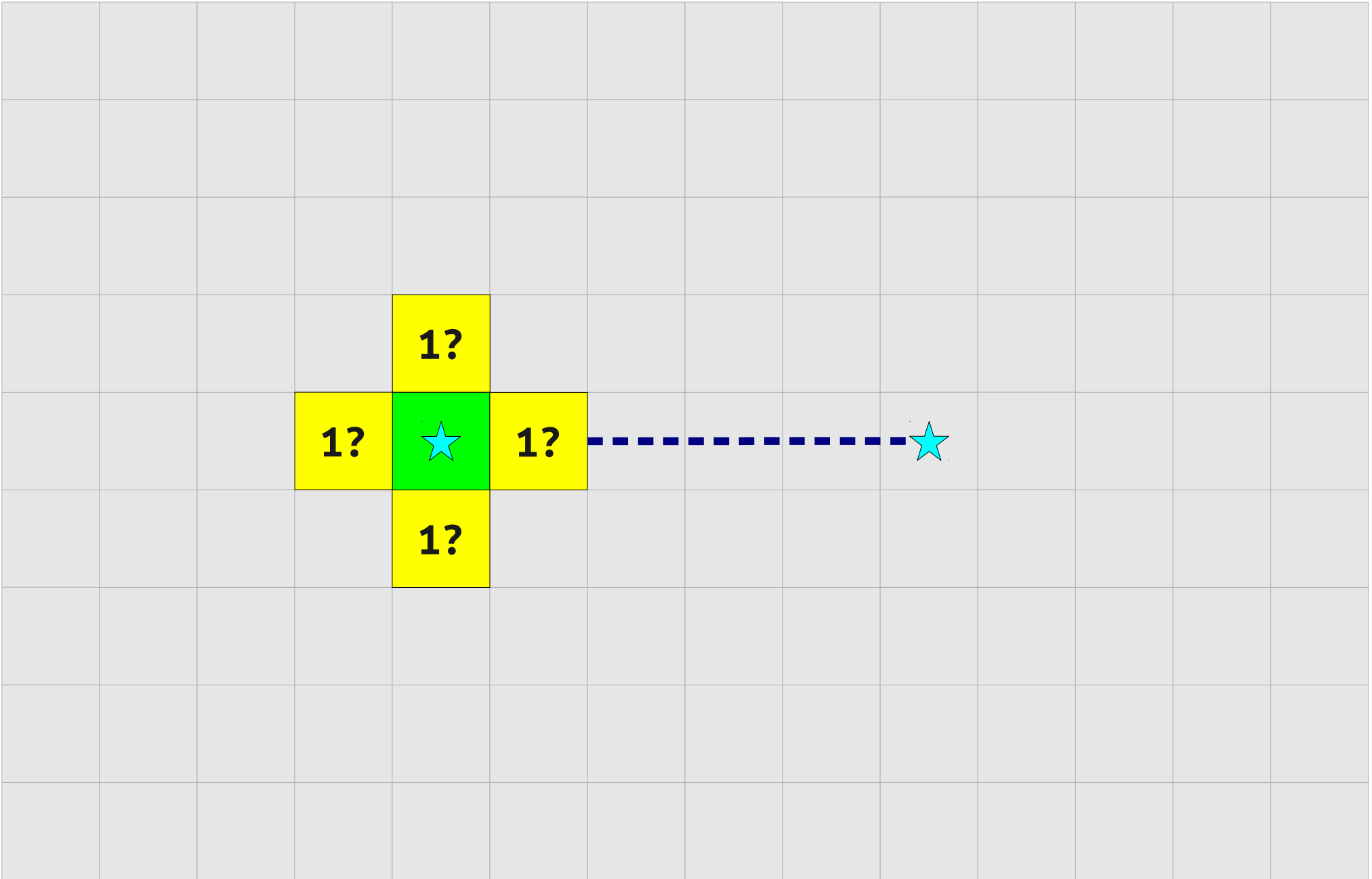


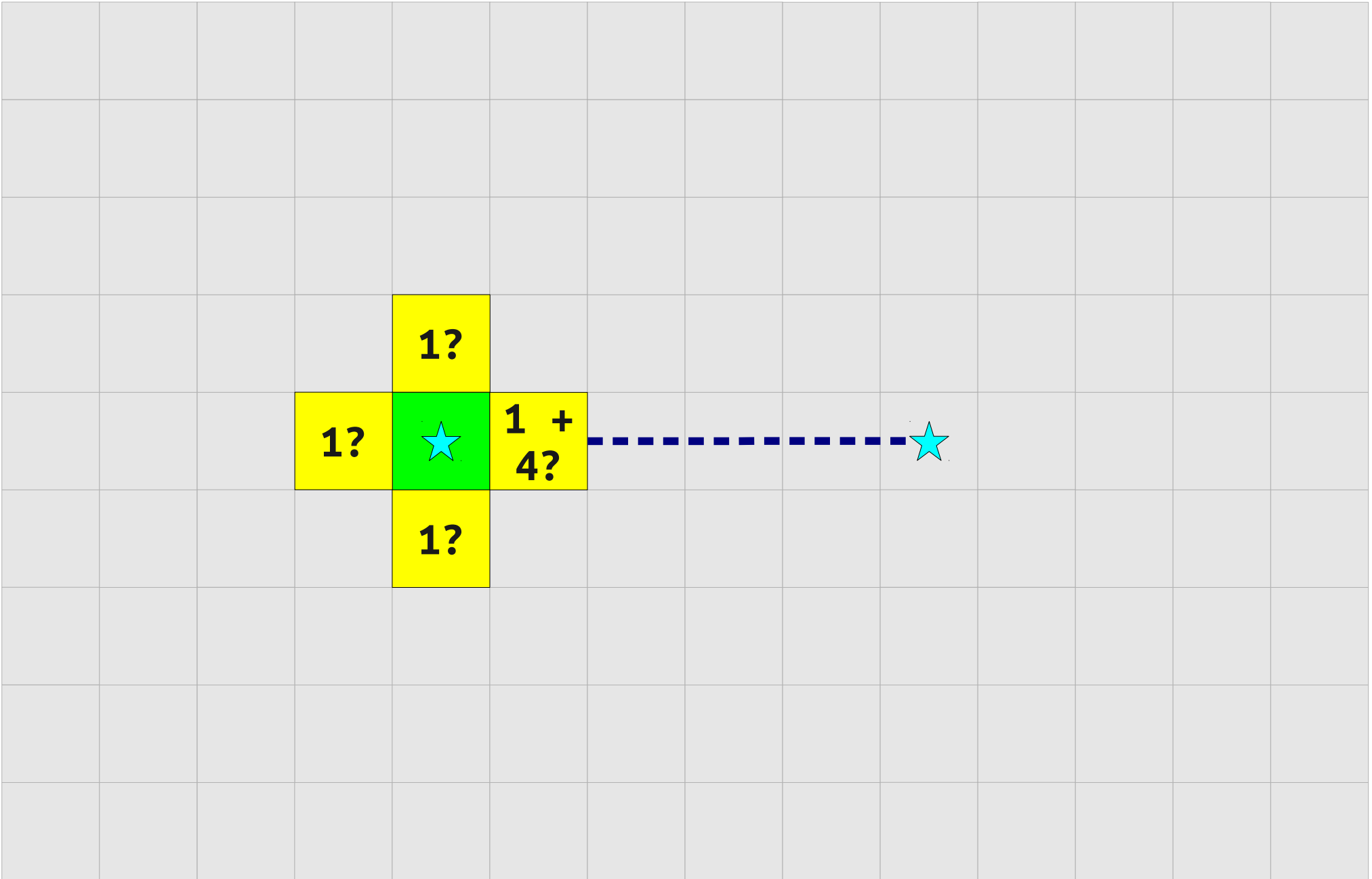


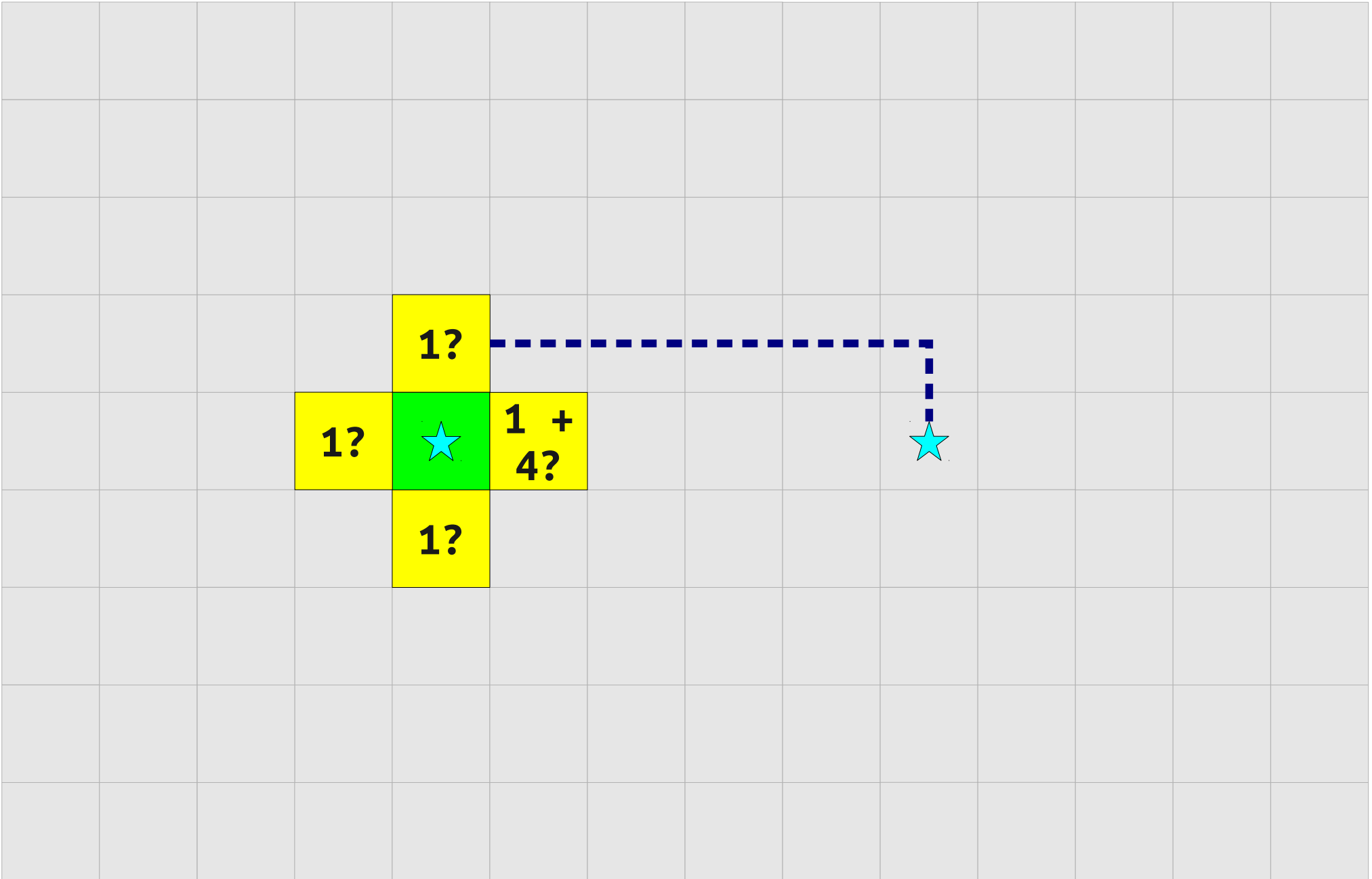


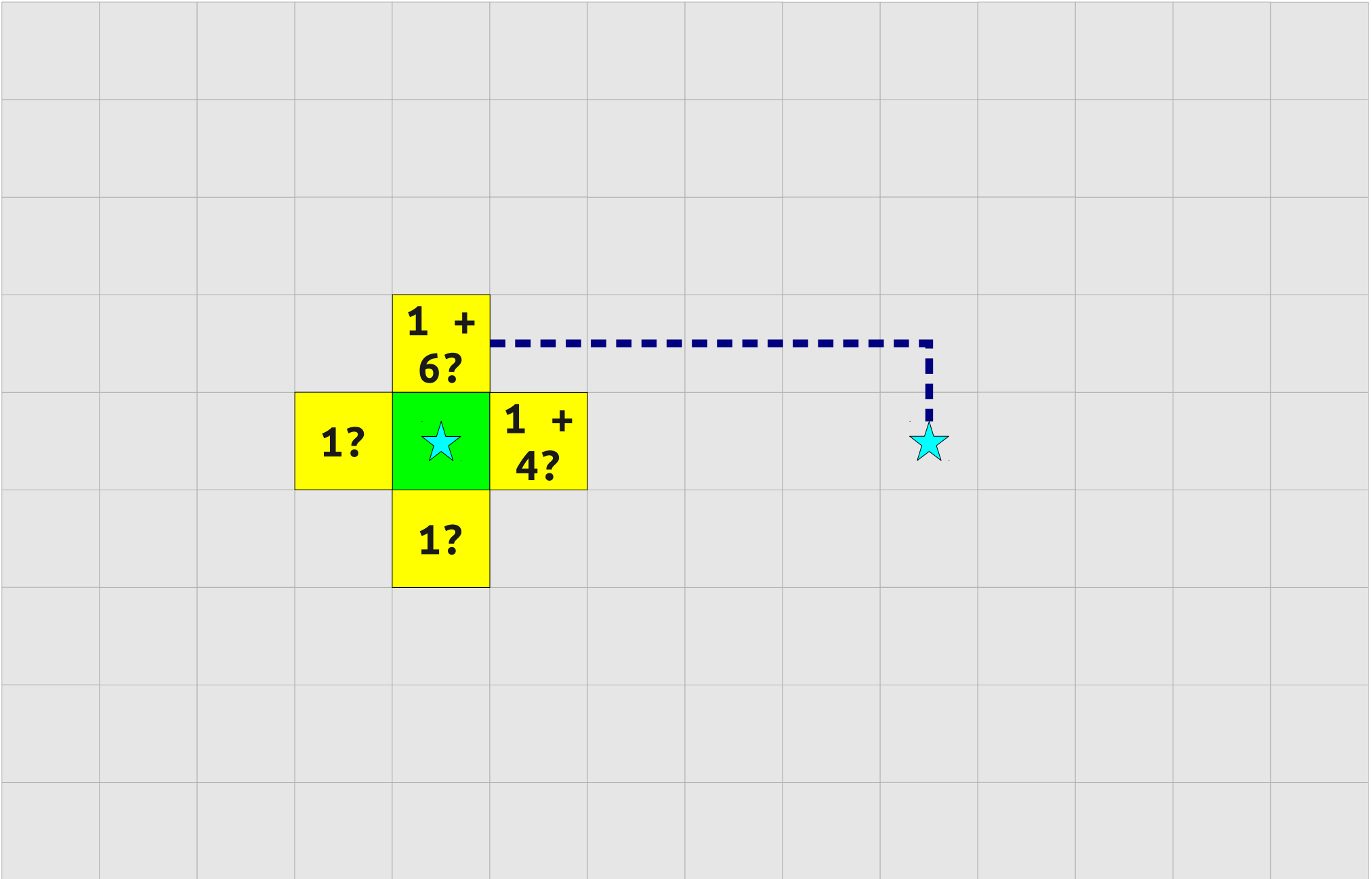


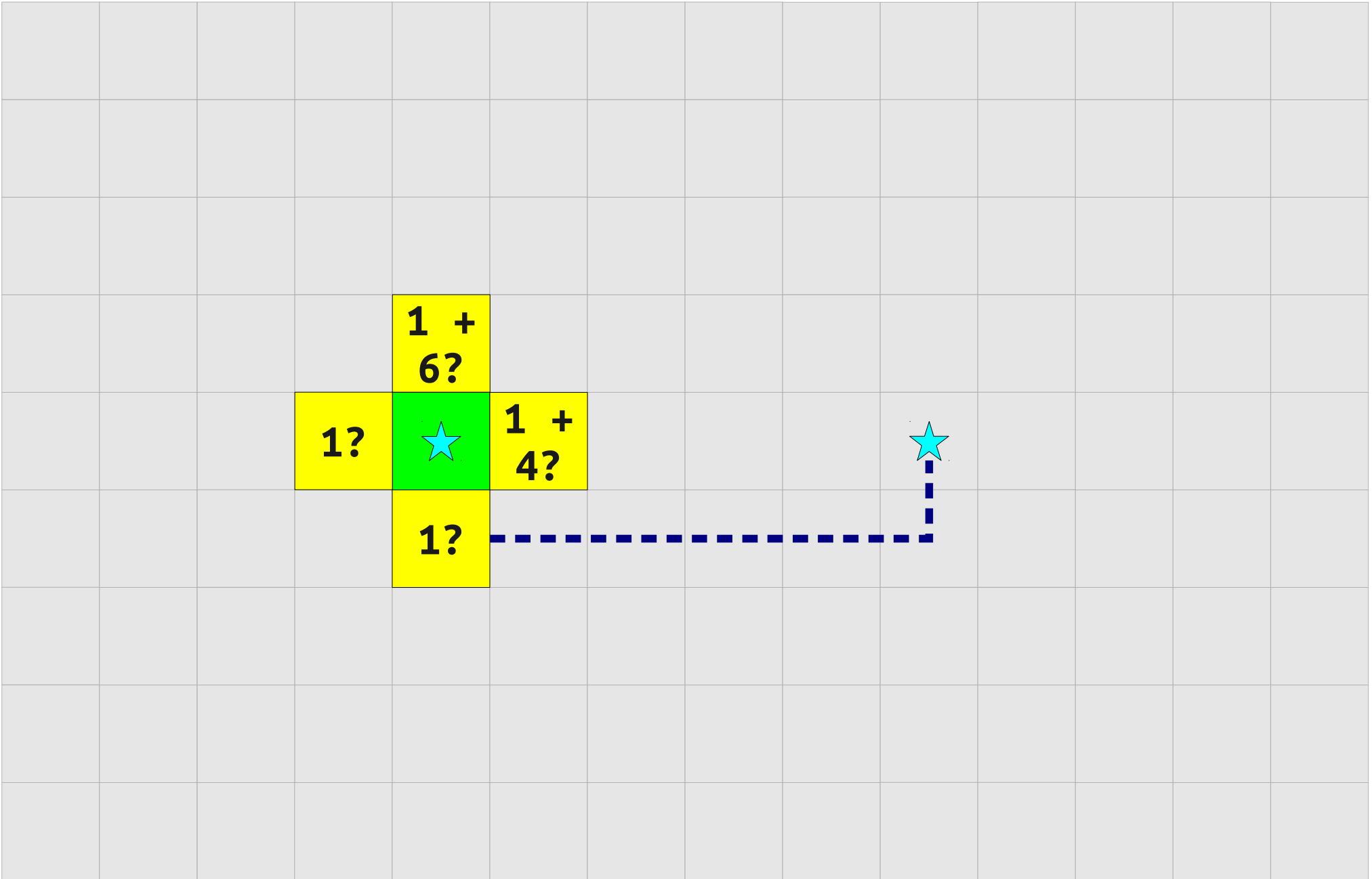


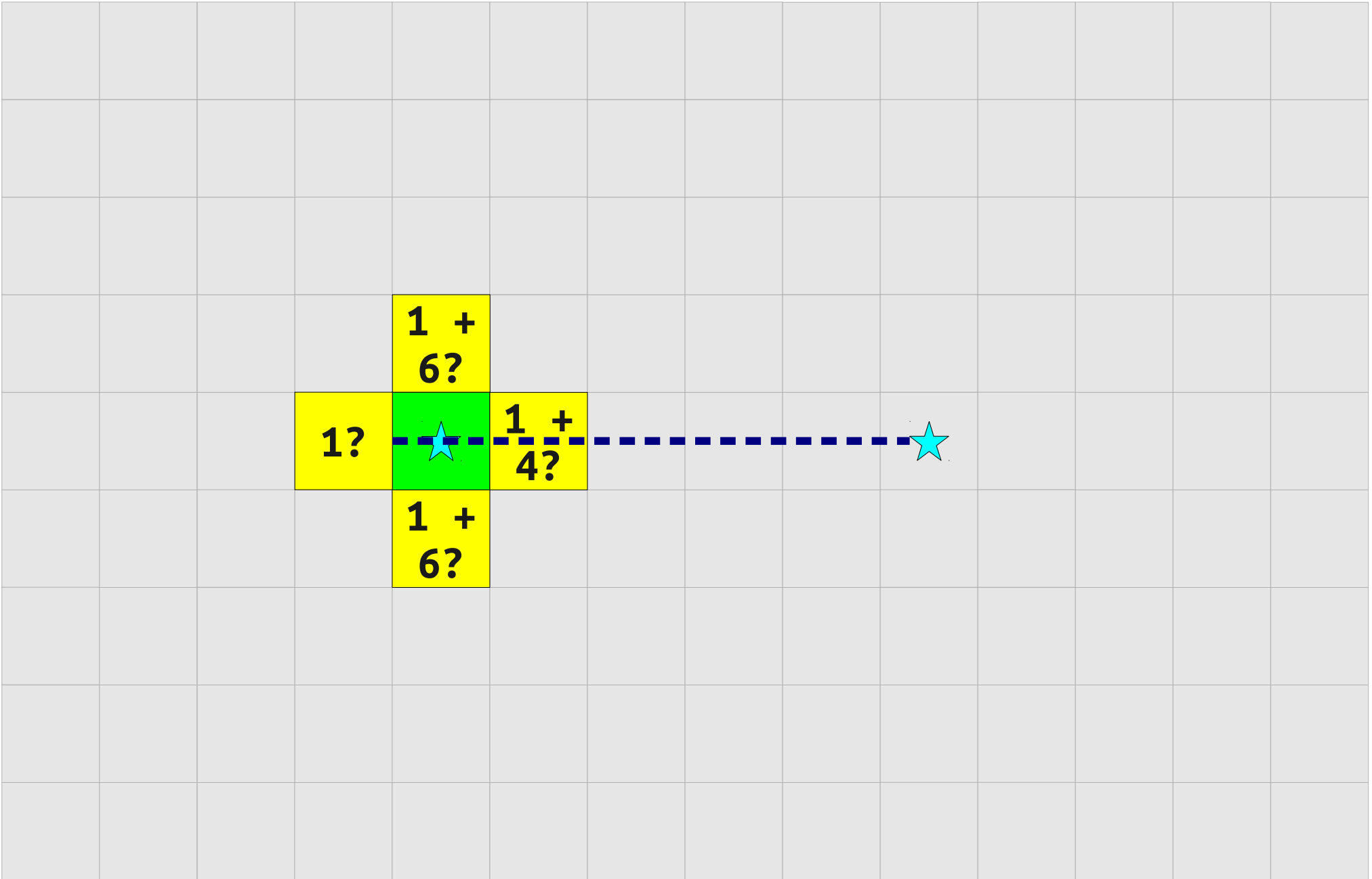


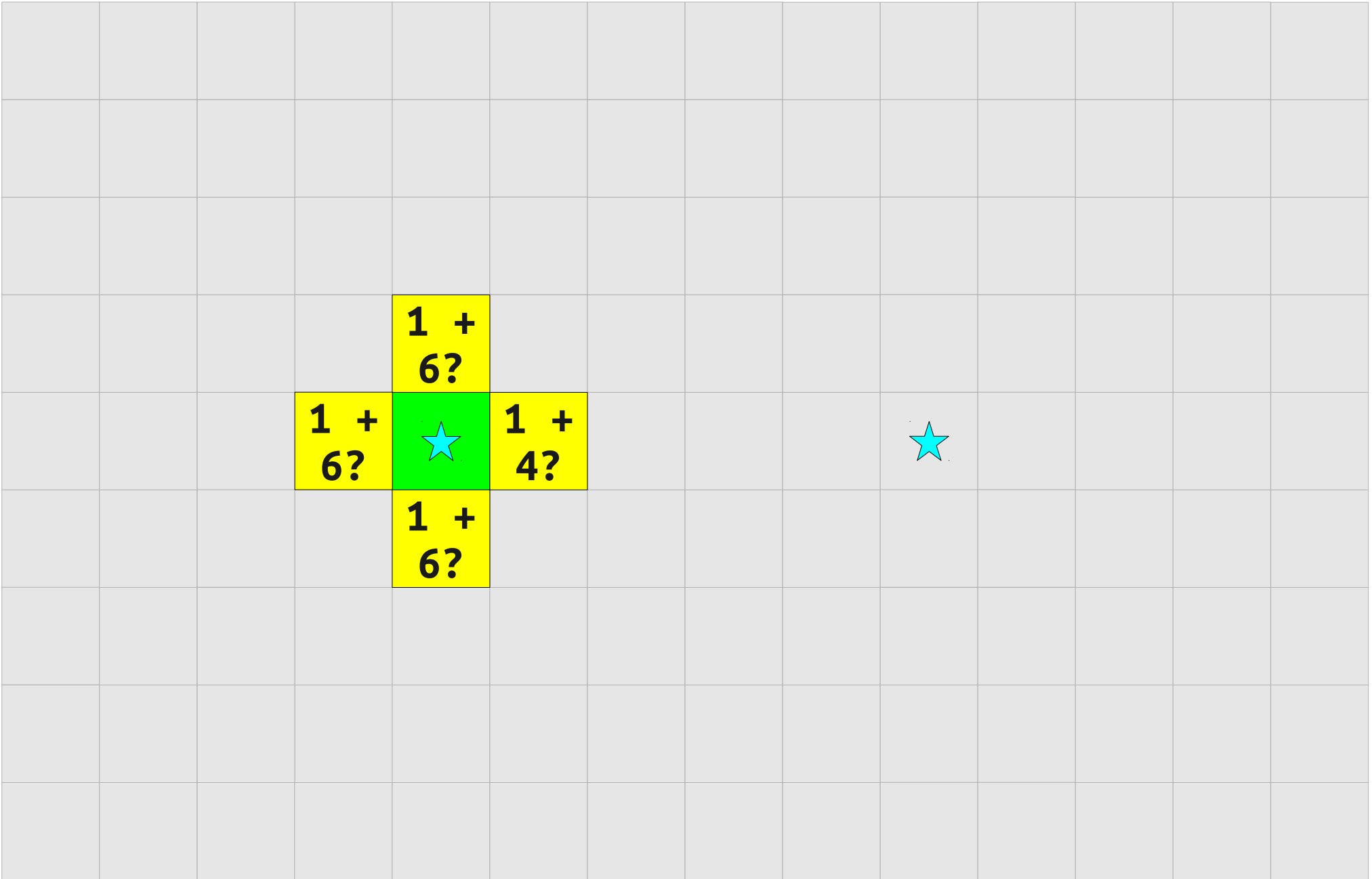












1 +
6?

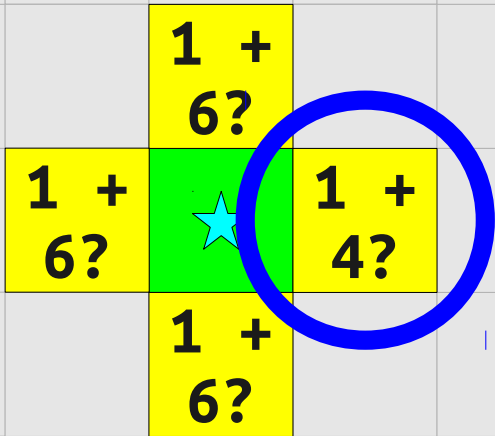
1 +
6?

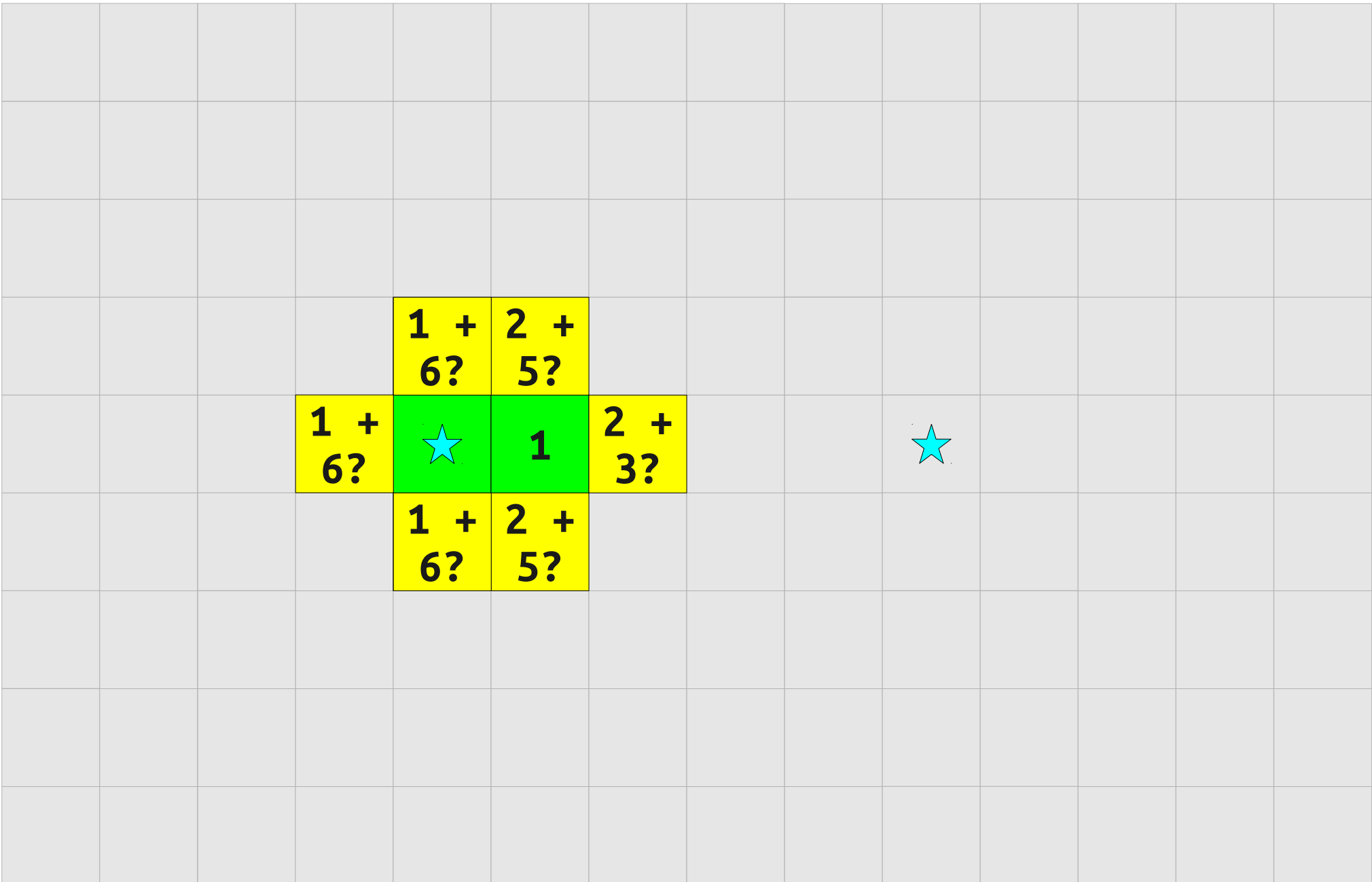


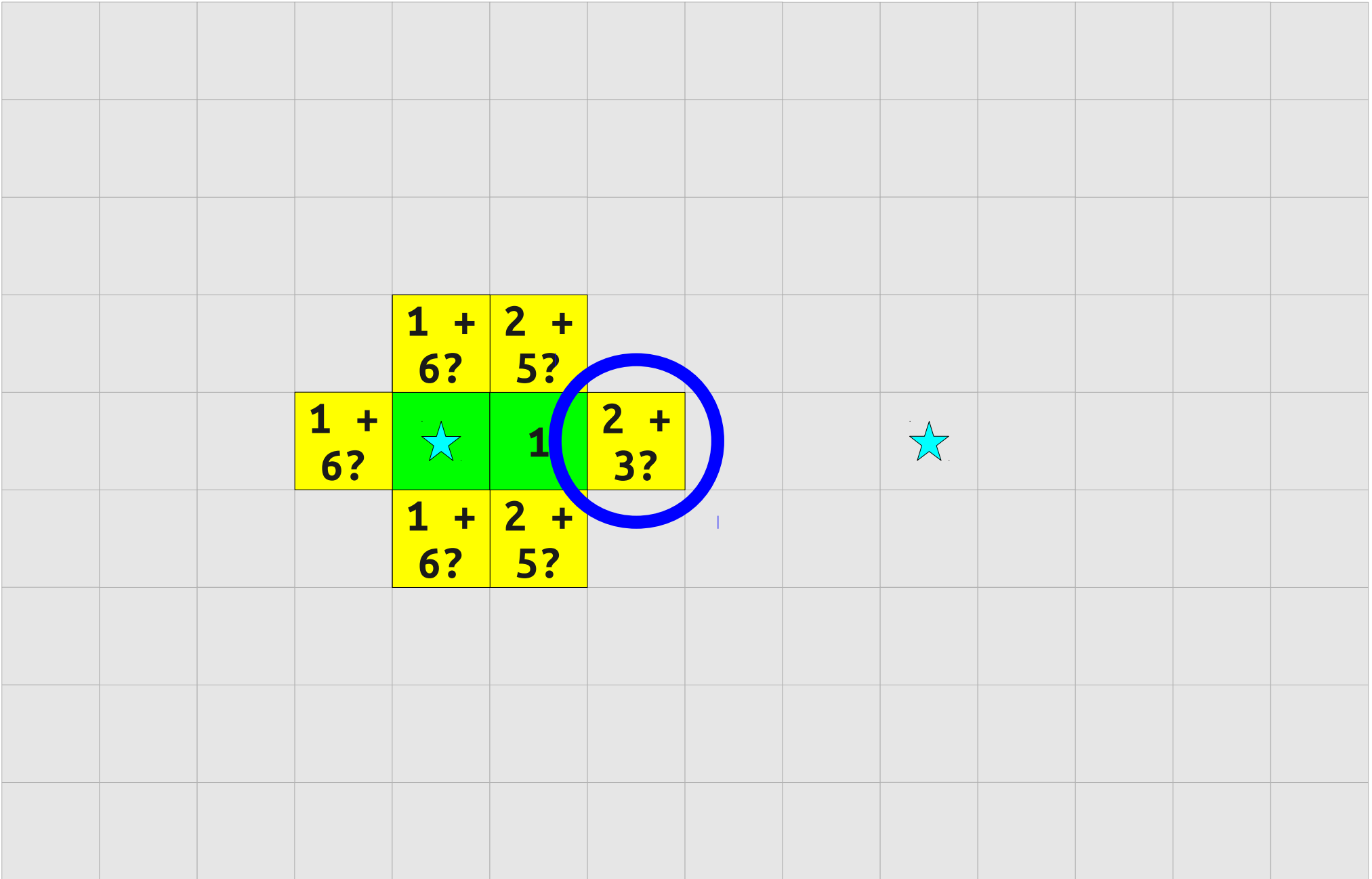
1 +
4?

1 +
6?

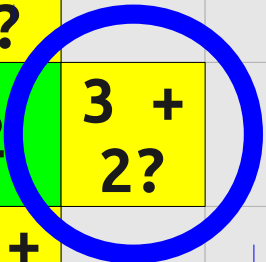






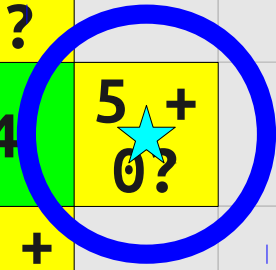


				1 + 6?	2 + 5?	3 + 4?								
			1 + 6?	★	1	2	3 + 2?							
				1 + 6?	2 + 5?	3 + 4?								



				1 + 6?	2 + 5?	3 + 4?	4 + 3?							
			1 + 6?	★	1	2	3	4 + 1?	★					
				1 + 6?	2 + 5?	3 + 4?	4 + 3?							

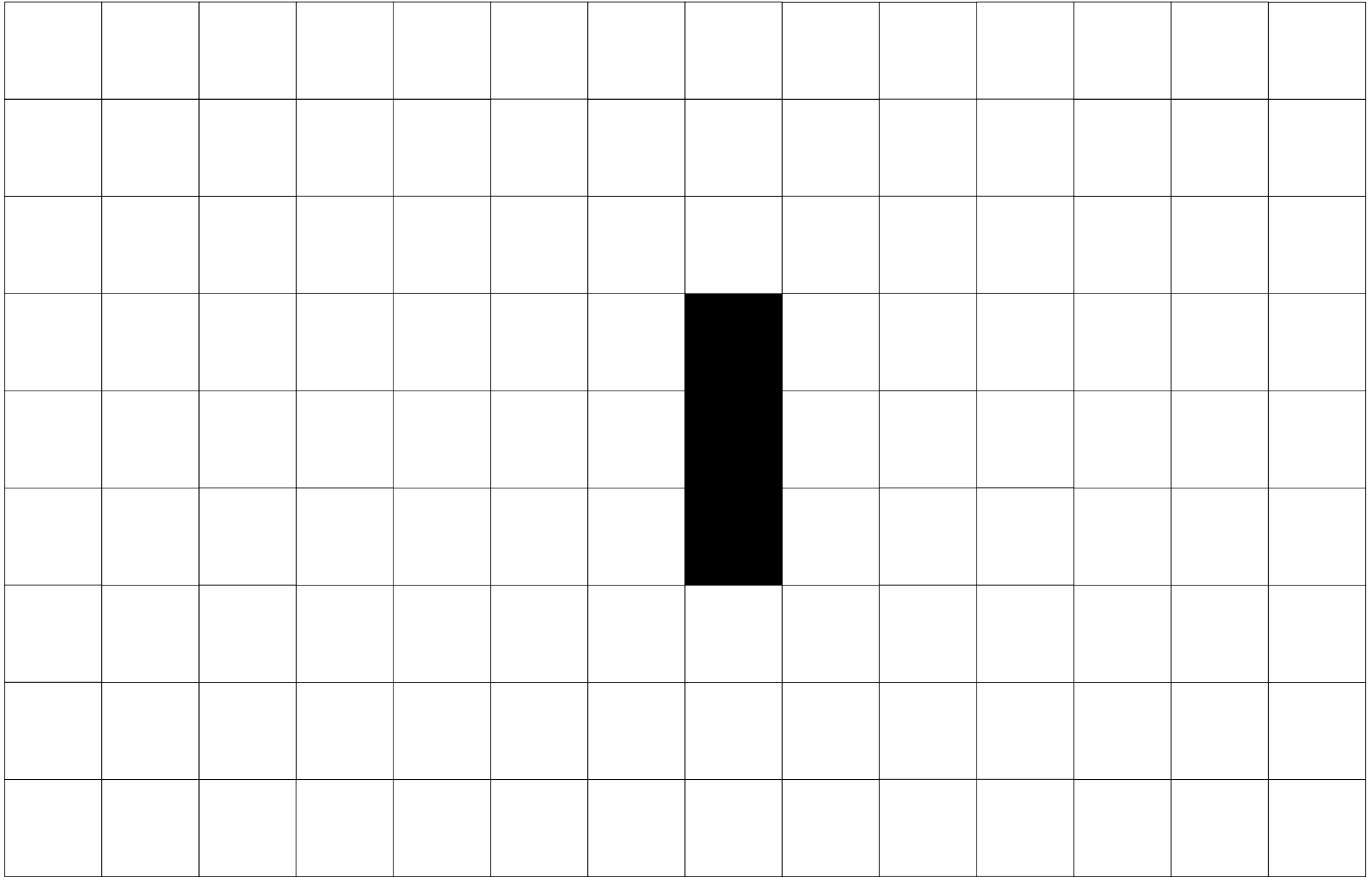
				1 + 6?	2 + 5?	3 + 4?	4 + 3?	5 + 2?					
			1 + 6?	★	1	2	3	4	5 + 6?				
				1 + 6?	2 + 5?	3 + 4?	4 + 3?	5 + 2?					

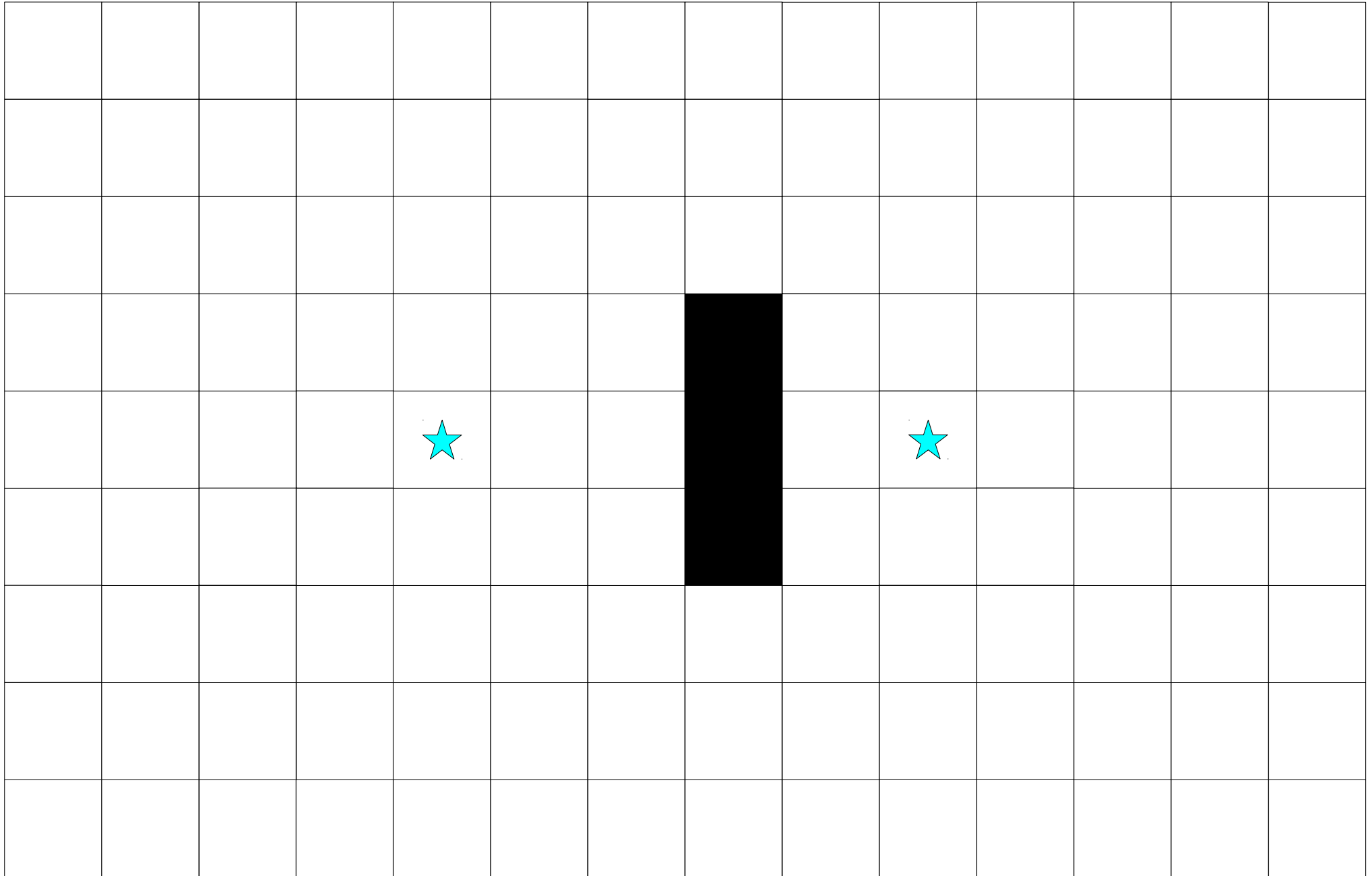


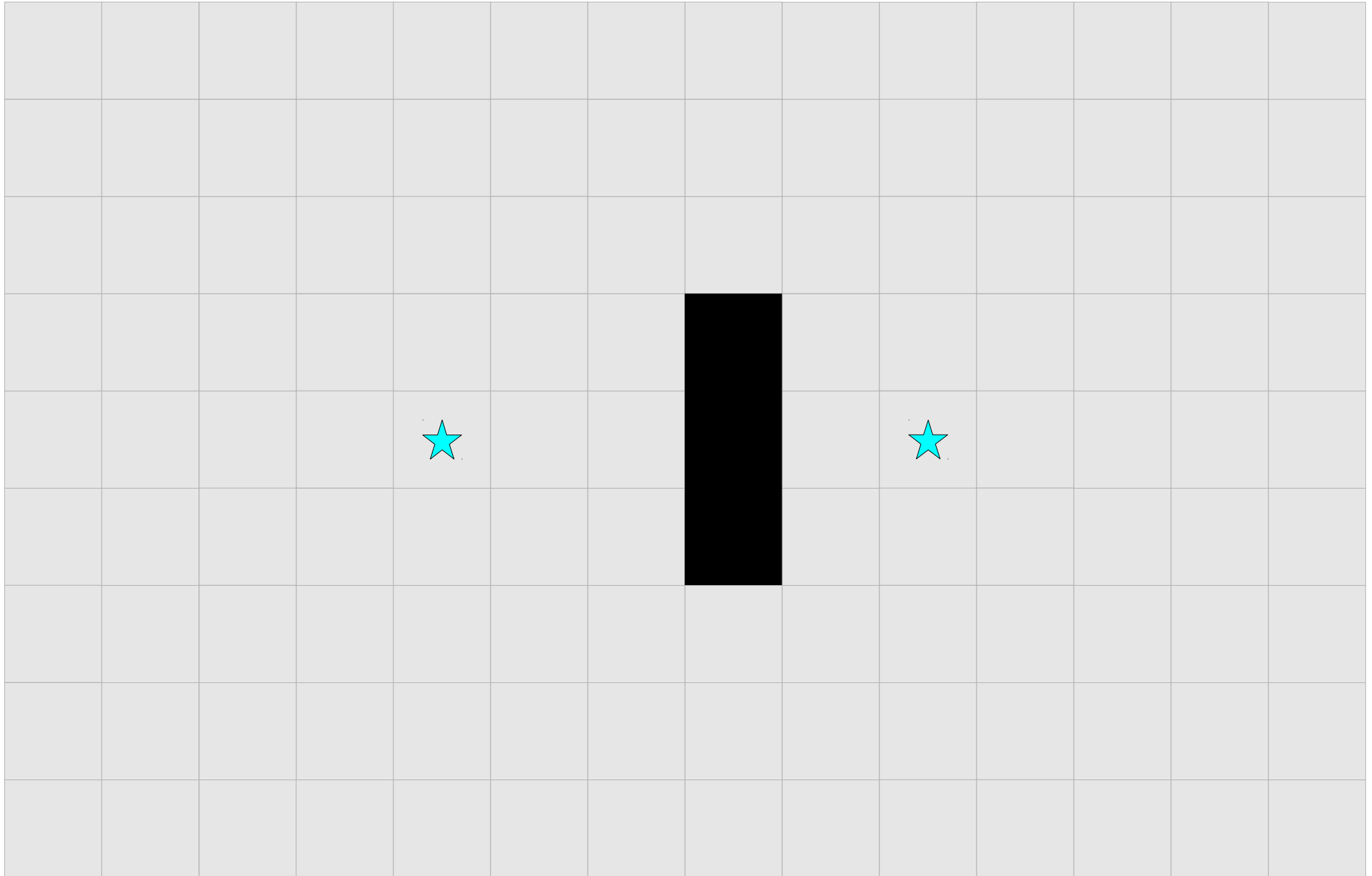
1 +	2 +	3 +	4 +	5 +
6?	5?	4?	3?	2?

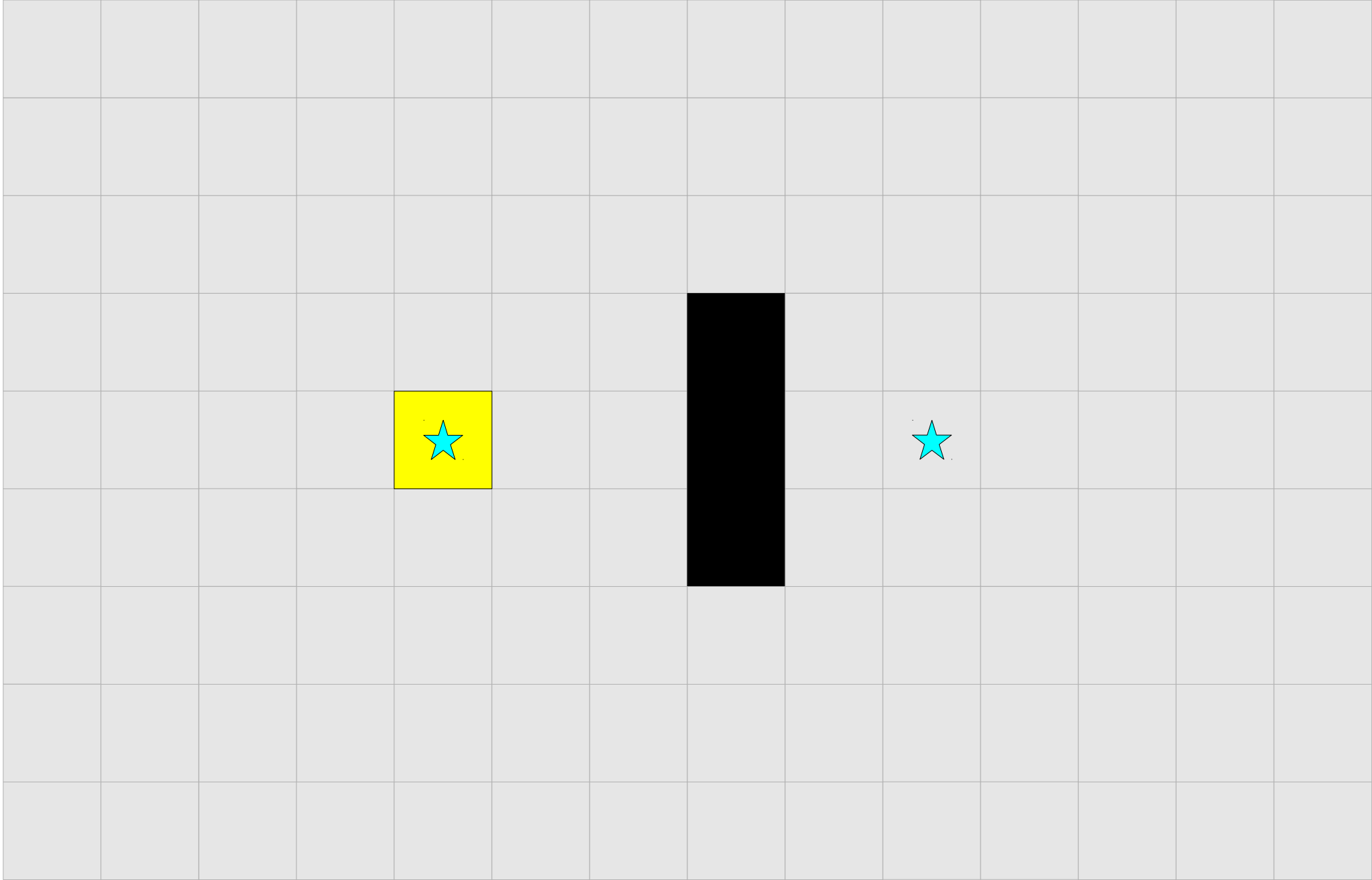
1 +	★	1	2	3	4	★
6?						

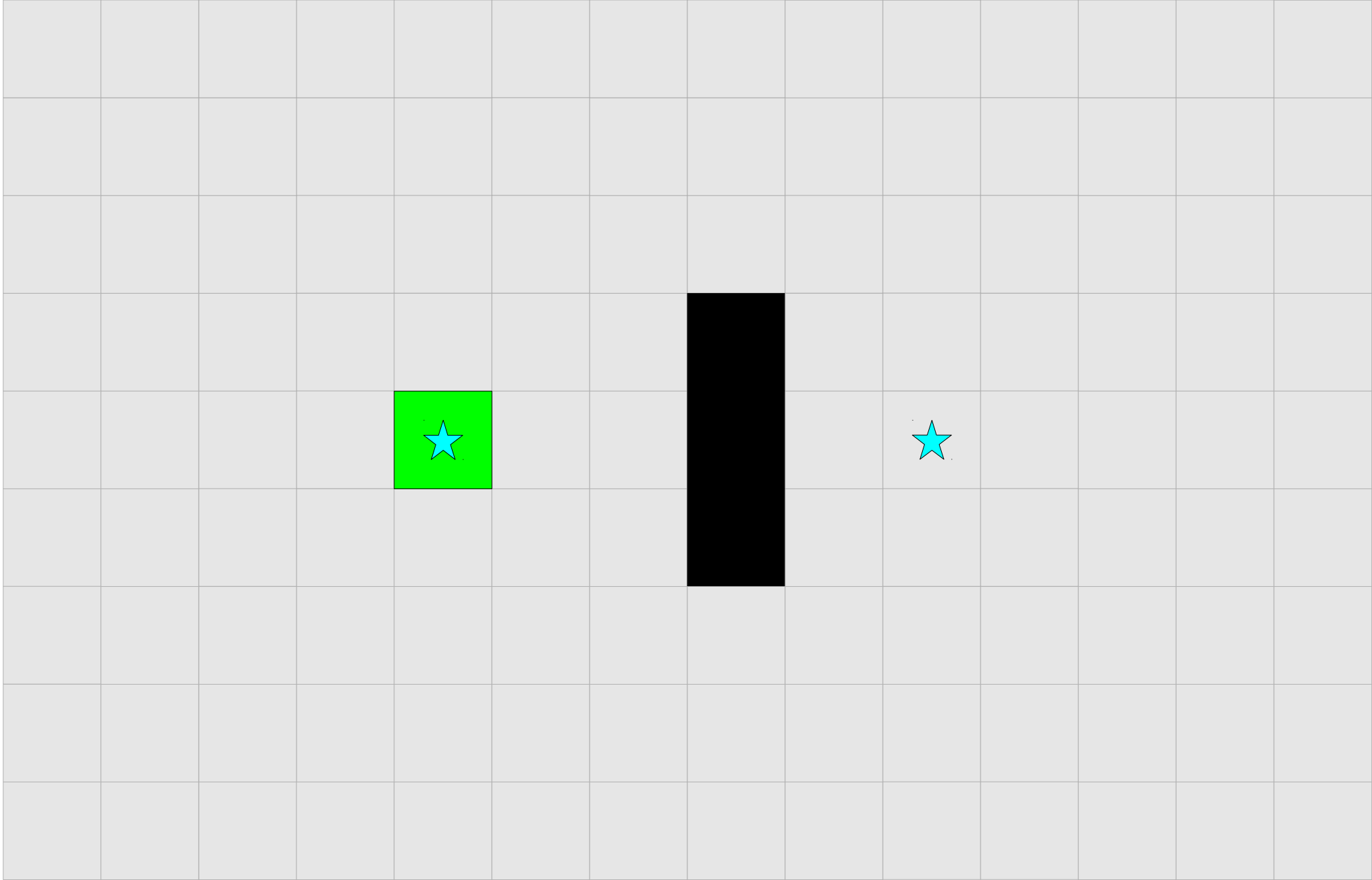
1 +	2 +	3 +	4 +	5 +
6?	5?	4?	3?	2?

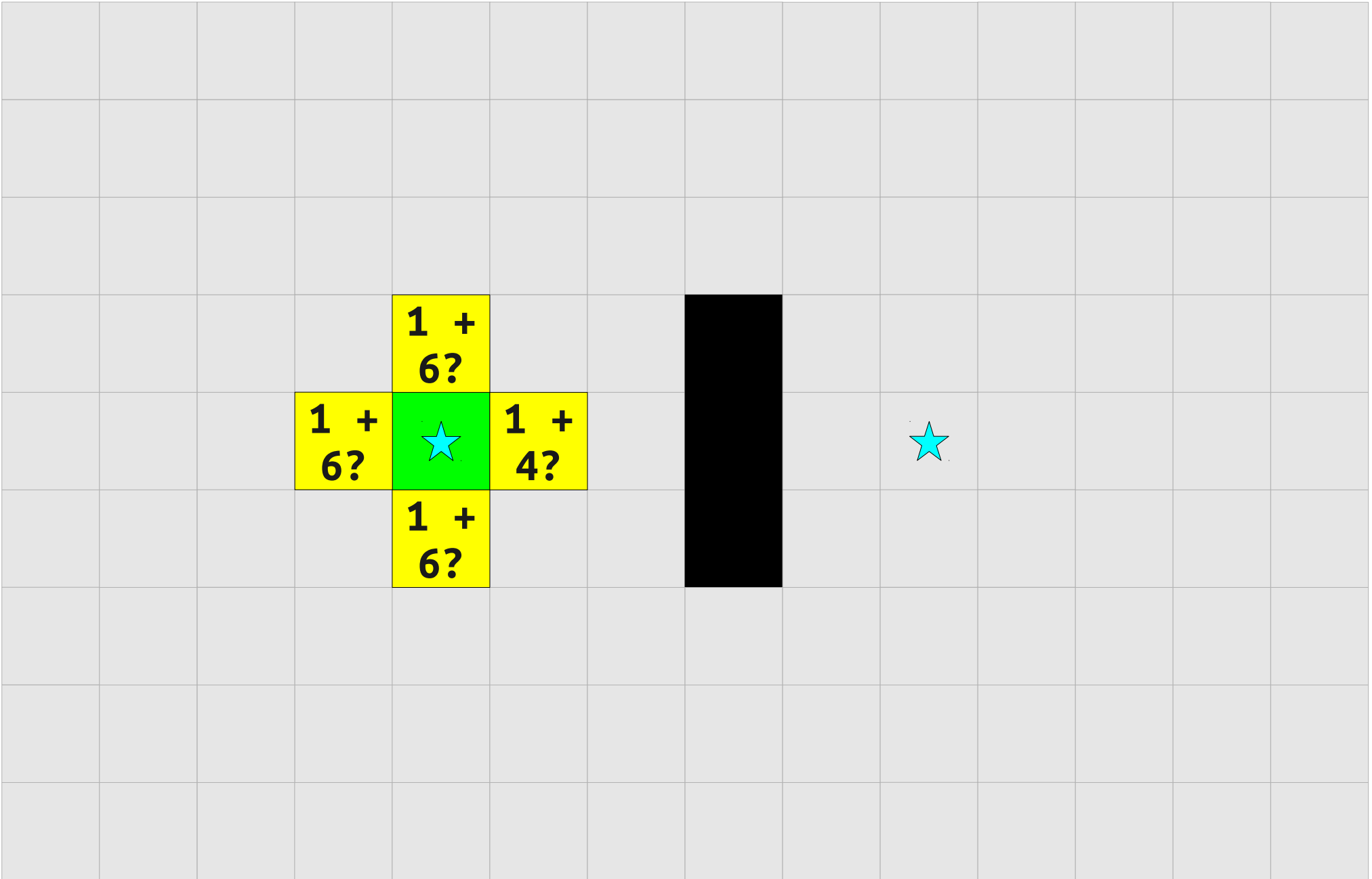


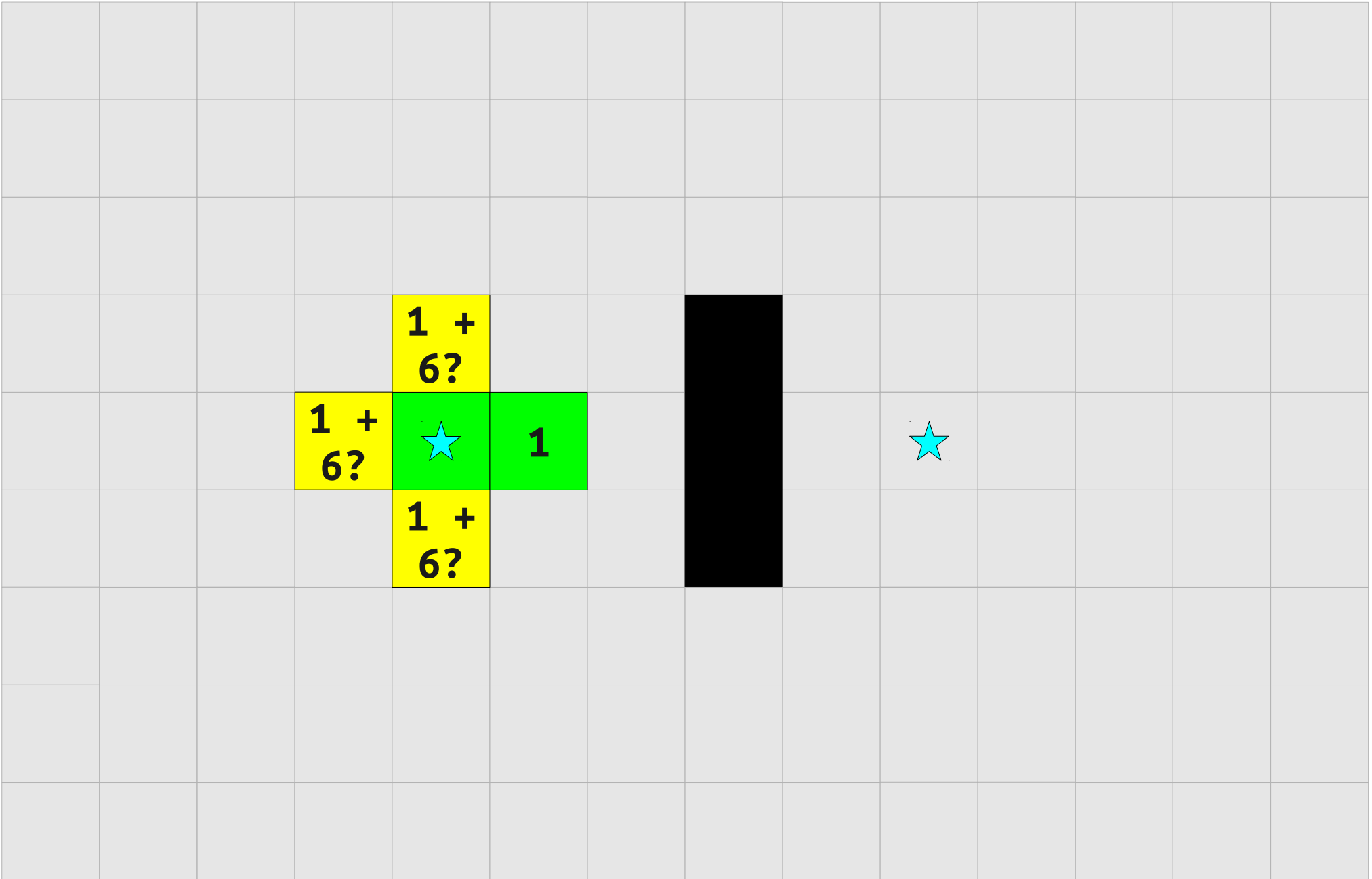












1 +
6?

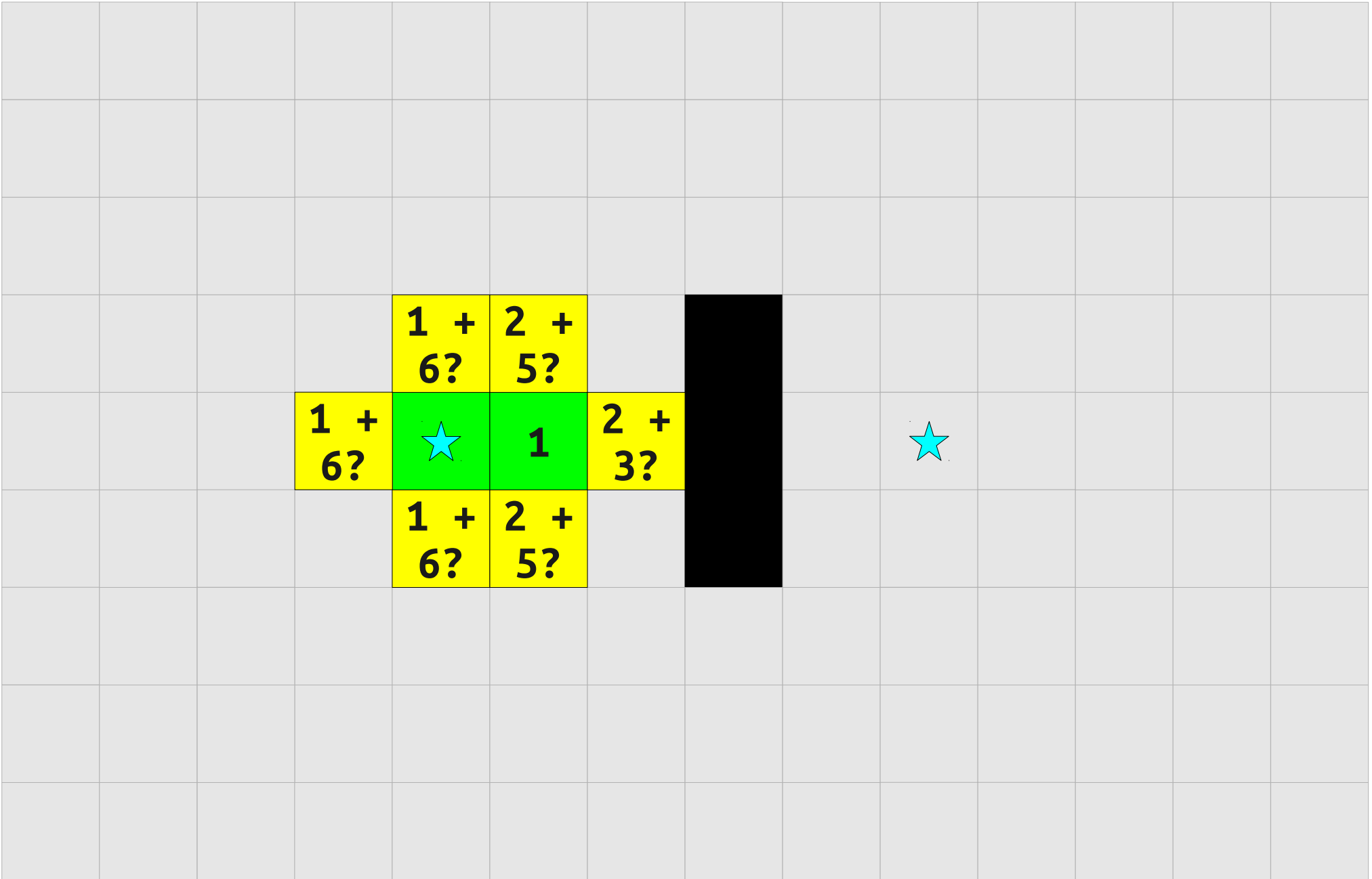
1 +
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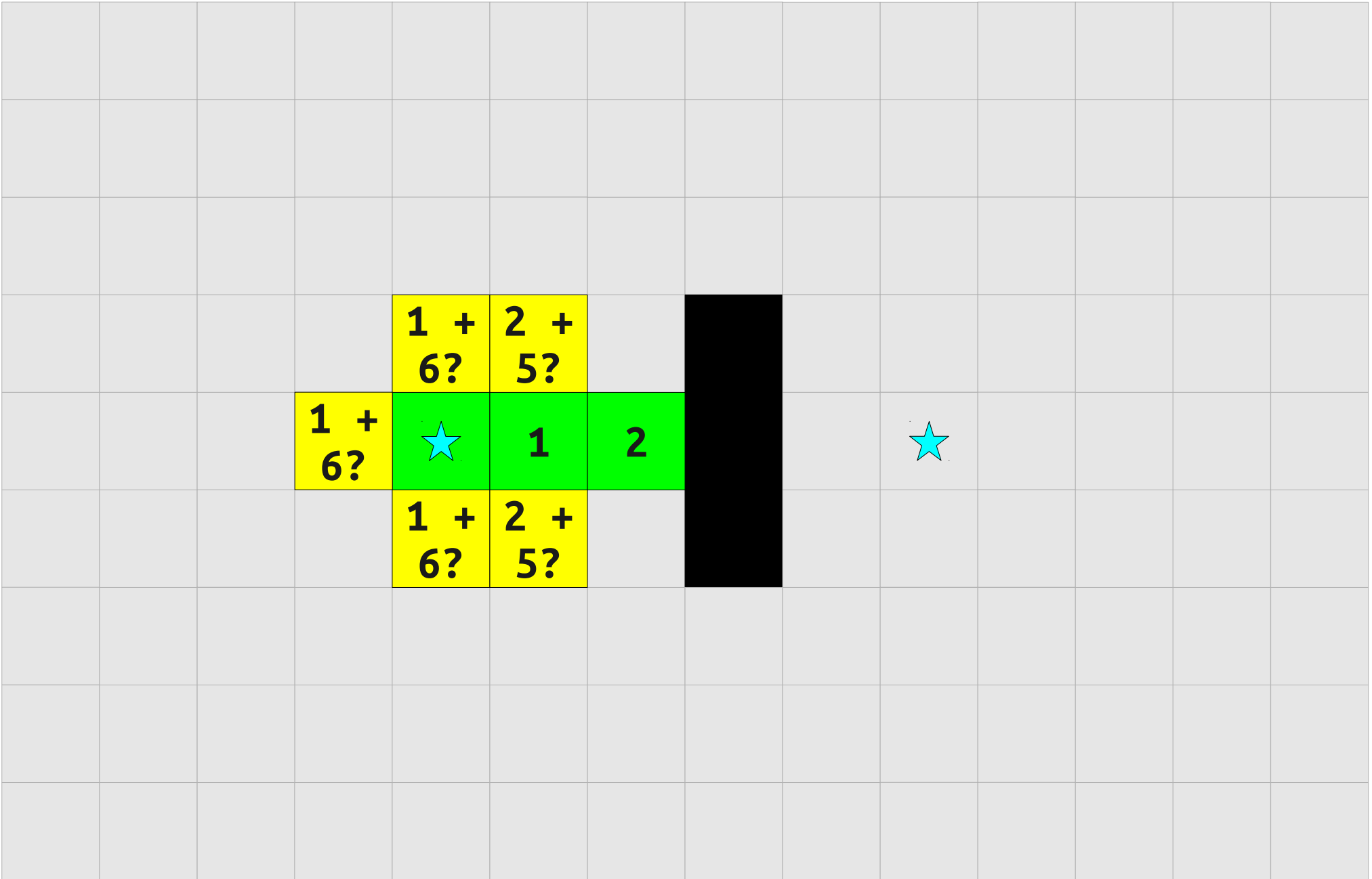


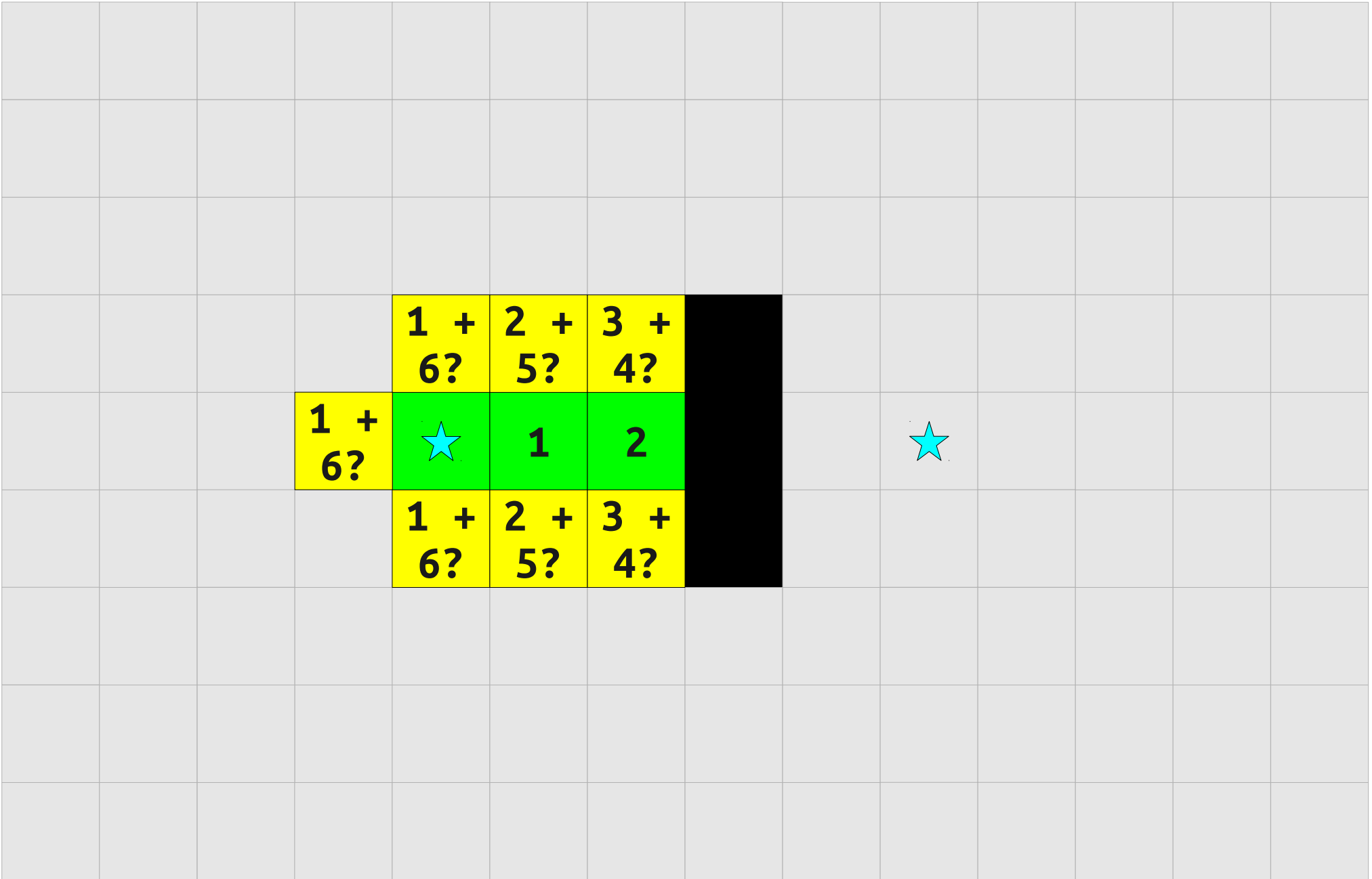
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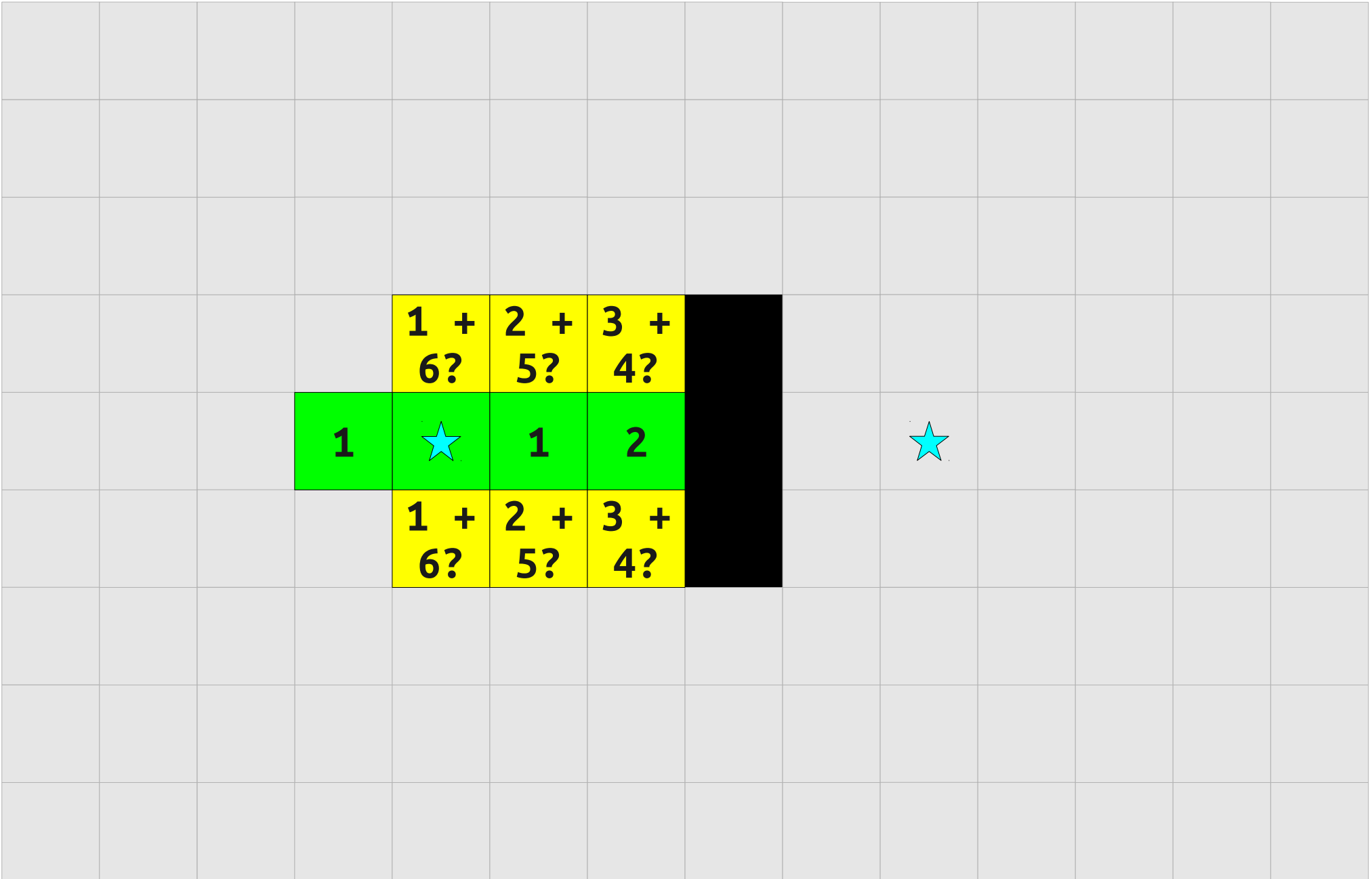
1 +
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			2 + 7?	1 + 6?	2 + 5?	3 + 4?							
		2 + 7?	1	★	1	2							
			2 + 7?	1 + 6?	2 + 5?	3 + 4?							



			$2 + 7?$	$1 + 6?$	$2 + 5?$	$3 + 4?$							
		$2 + 7?$	1	★	1	2							
			$2 + 7?$	$1 + 6?$	2	$3 + 4?$							

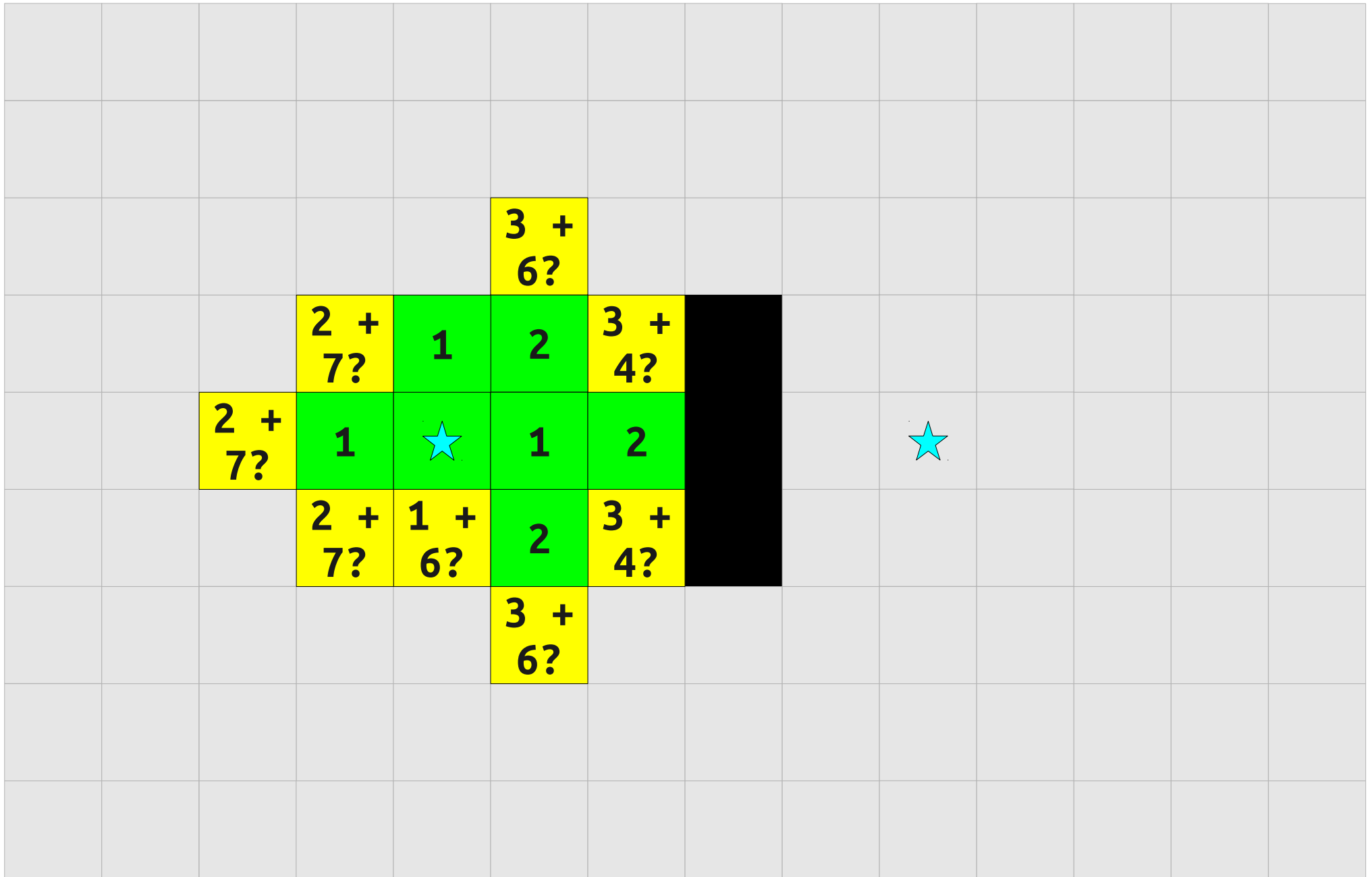


			$2 + 7?$	$1 + 6?$	$2 + 5?$	$3 + 4?$							
		$2 + 7?$	1	★	1	2							
			$2 + 7?$	$1 + 6?$	2	$3 + 4?$							
					$3 + 6?$								



			2 + 7?	1 + 6?	2	3 + 4?	
		2 + 7?	1	★	1	2	
			2 + 7?	1 + 6?	2	3 + 4?	
					3 + 6?		





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7? 6?

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7? 6? 2 4?

3 +
6?



2 + 3 +
7? 6?

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2 + 3 +
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2 + 3 + 4 +
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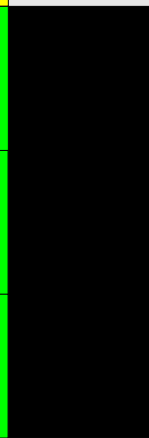
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$$\begin{array}{r} 3 + \\ 8? \end{array}$$

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$$\begin{array}{r} 4 + \\ 7? \end{array}$$

$$\begin{array}{r} 5 + \\ 6? \end{array}$$

For Comparison: What Dijkstra's Algorithm Would Have Searched



A* Search

- As long as the heuristic is admissible (and satisfies one other technical condition), A* will always find the shortest path from the source to the destination node.
- Can be *dramatically* faster than Dijkstra's algorithm.
 - Focuses work in areas likely to be productive.
 - Avoids solutions that appear worse until there is evidence they may be appropriate.

Close Cousins

- Dijkstra's algorithm and A* search are very closely related:
 - Dijkstra's uses a node's candidate distance as its priority.
 - A* uses a node's candidate distance **plus a heuristic value** as its priority.
- Interesting fact: If you use the **zero heuristic** (which always predicts a node is at distance 0 from the endpoint), A* search is completely identical to Dijkstra's algorithm!

Dijkstra's Algorithm

- Mark all nodes as gray.
- Mark the initial node s as yellow and at candidate distance 0 .
- Enqueue s into the priority queue with priority 0 .
- While not all nodes have been visited:
 - Dequeue the lowest-cost node u from the priority queue.
 - Color u green. The candidate distance d that is currently stored for node u is the length of the shortest path from s to u .
 - If u is the destination node t , you have found the shortest path from s to t and are done.
 - For each node v connected to u by an edge of length L :
 - If v is gray:
 - Color v yellow.
 - Mark v 's distance as $d + L$.
 - Set v 's parent to be u .
 - Enqueue v into the priority queue with priority $d + L$.
 - If v is yellow and the candidate distance to v is greater than $d + L$:
 - Update v 's candidate distance to be $d + L$.
 - Update v 's parent to be u .
 - Update v 's priority in the priority queue to $d + L$.

A* Search

- Mark all nodes as gray.
- Mark the initial node s as yellow and at candidate distance 0 .
- Enqueue s into the priority queue with priority $h(s, t)$.
- While not all nodes have been visited:
 - Dequeue the lowest-cost node u from the priority queue.
 - Color u green. The candidate distance d that is currently stored for node u is the length of the shortest path from s to u .
 - If u is the destination node t , you have found the shortest path from s to t and are done.
 - For each node v connected to u by an edge of length L :
 - If v is gray:
 - Color v yellow.
 - Mark v 's distance as $d + L$.
 - Set v 's parent to be u .
 - Enqueue v into the priority queue with priority $d + L + h(v, t)$.
 - If v is yellow and the candidate distance to v is greater than $d + L$:
 - Update v 's candidate distance to be $d + L$.
 - Update v 's parent to be u .
 - Update v 's priority in the priority queue to $d + L + h(v, t)$.

Next Time

- **More Graph Algorithms**
 - Minimum Spanning Trees
 - Kruskal's Algorithm
 - Data Clustering