

CS 106B, Lecture 18

Binary Search Trees

Plan for Today

- Start with a discussion of how to implement a Set
 - The importance of choosing a good data structure
- Move into trees, a new kind of data structure
- We'll focus on "reading" trees today – modifying trees will be tomorrow's lecture

Designing a Set

- We've seen how to implement:

- Stack (array or linked list)
 - Vector (array)
 - Queue (linked list)

- How would we implement Set?

- Add
 - Contains
 - Remove

First Try

- Store all the elements in an **unsorted** array or linked list
 - What is the Big-Oh of contains?
 - What is the Big-Oh of adding an element?
 - What is the Big-Oh of removing an element?

0	1	2	3	4	5	6	7	8	9	10
3	8	9	7	5	12	4	8	1	6	75

Another attempt

- What if we **sorted** the array?
 - What is the Big-Oh of contains?
 - What is the Big-Oh of adding an element?
 - What is the Big-Oh of removing an element?

0	1	2	3	4	5	6	7	8	9	10
2	5	6	8	11	13	17	22	23	29	31

Binary Search

- Fast way to search for elements in a **sorted array**
- Looping through elements one by one is slow [$O(N)$]
- Idea:

Jump to the middle element:

- if the middle is what we're looking for, we're done. Hooray!
- if the middle is too small (we didn't go far enough) – we rule out the entire **left side** of elements smaller than the middle element
- if the middle is too big (we went too far) – we rule out the entire **right side** of elements bigger than the middle element

0	1	2	3	4	5	6	7	8	9	10
2	5	6	8	11	13	17	22	23	29	31

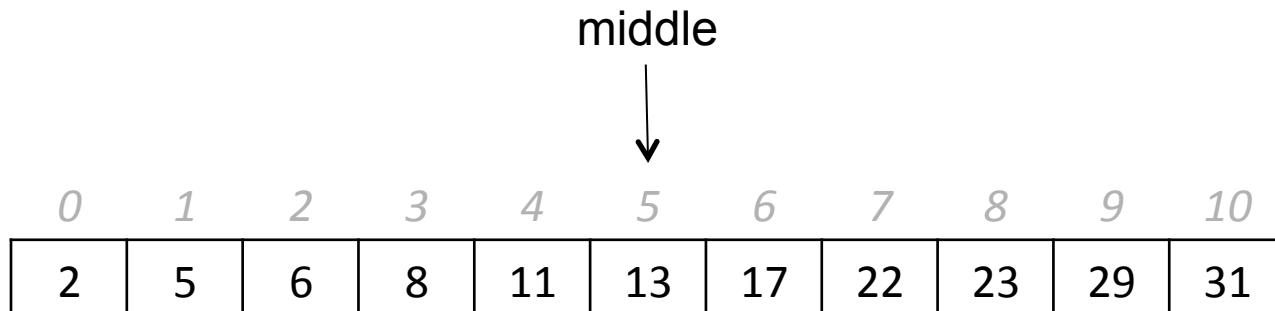
Binary Search in Action

- Search for 8:

0	1	2	3	4	5	6	7	8	9	10
2	5	6	8	11	13	17	22	23	29	31

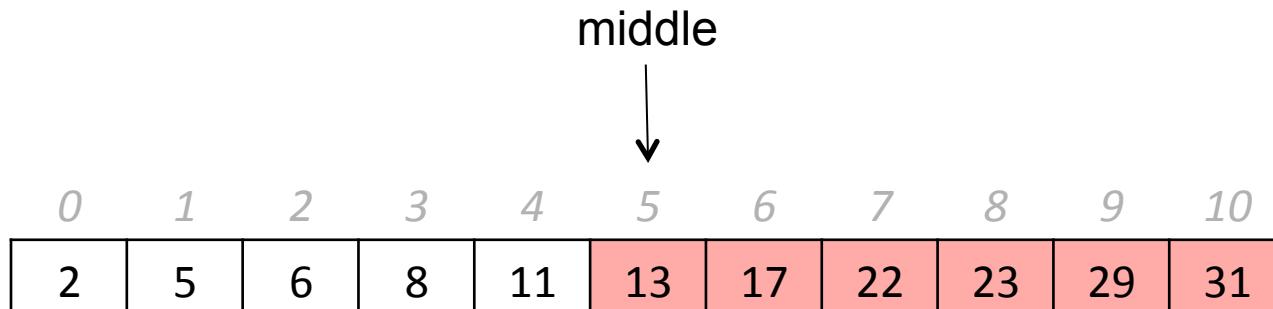
Binary Search in Action

- Search for 8:



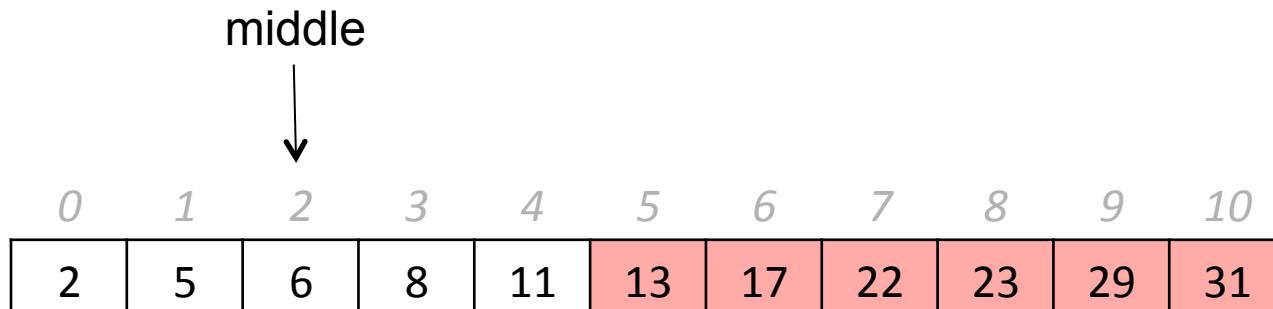
Binary Search in Action

- Search for 8:
- Look at 13
 - it's too big, so we rule out indices 5-10



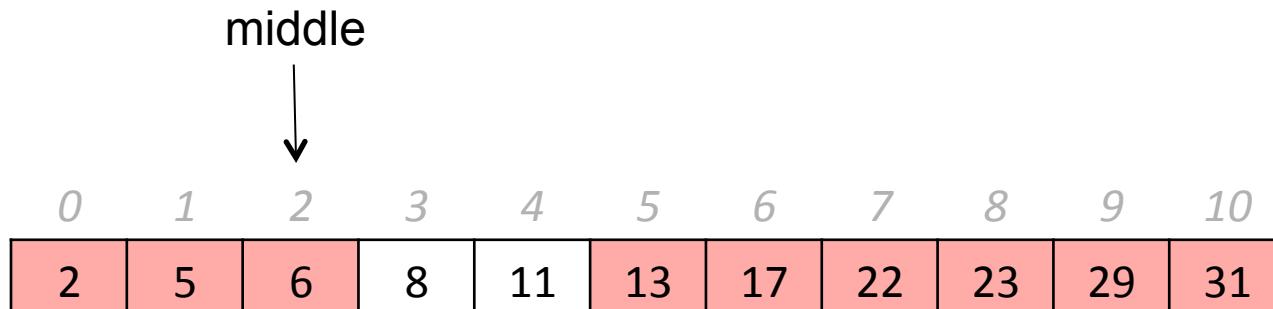
Binary Search in Action

- Search for 8:
- Look at 13
 - it's too big, so we rule out indices 5-10
- Pick the new middle of the remaining elements
- Look at 6:



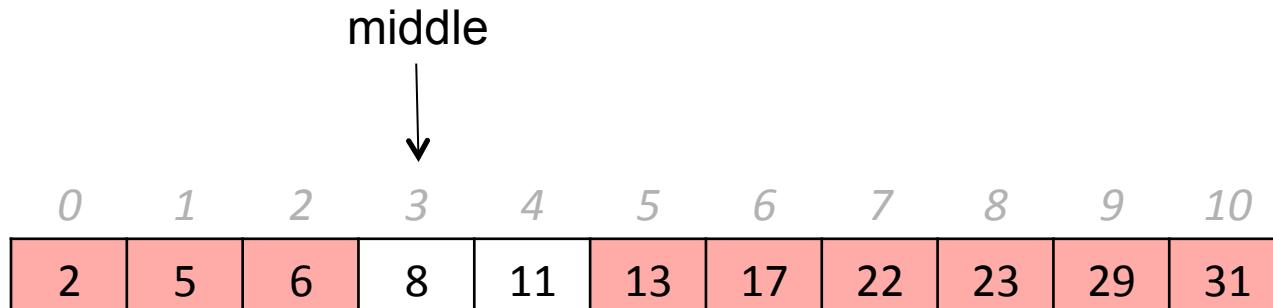
Binary Search in Action

- Search for 8:
- Look at 13
 - it's too big, so we rule out indices 5-10
- Pick the new middle of the remaining elements
- Look at 6:
 - it's too small, so we rule out indices 0-3



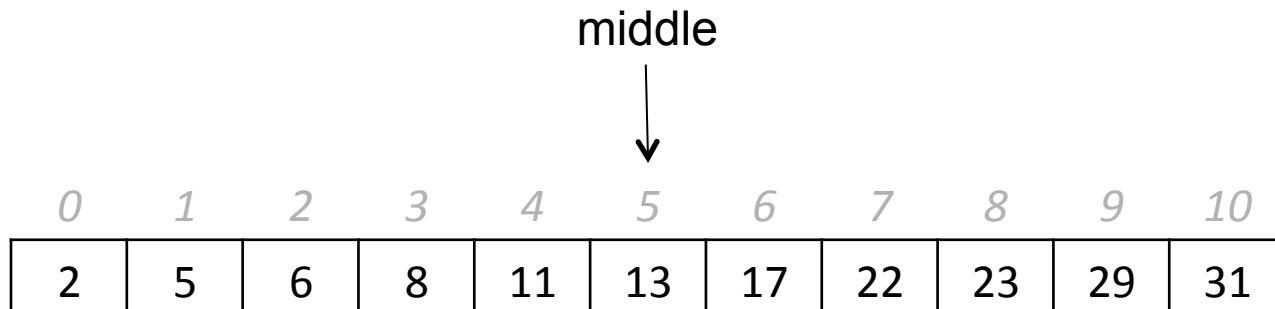
Binary Search in Action

- Search for 8:
- Look at 13
 - it's too big, so we rule out indices 5-10
- Pick the new middle of the remaining elements
- Look at 6:
 - it's too small, so we rule out indices 0-3
- Look at 8:
 - it's just right! We return true



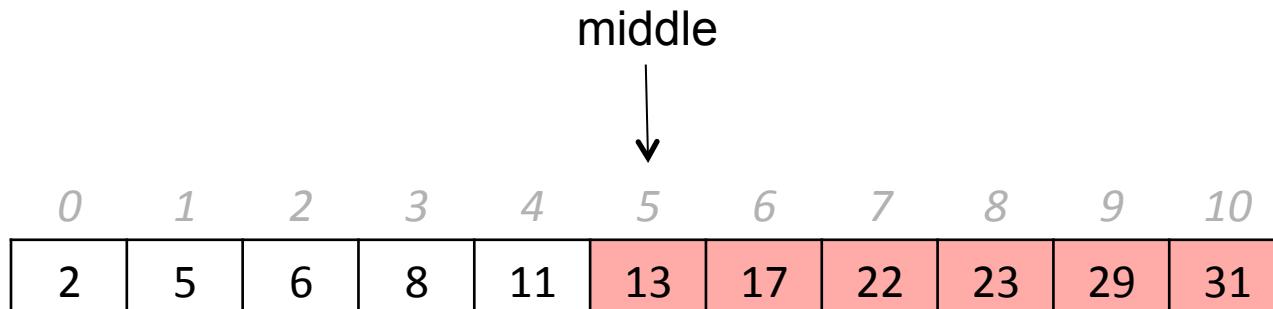
Binary Search in Action

- Search for 7:



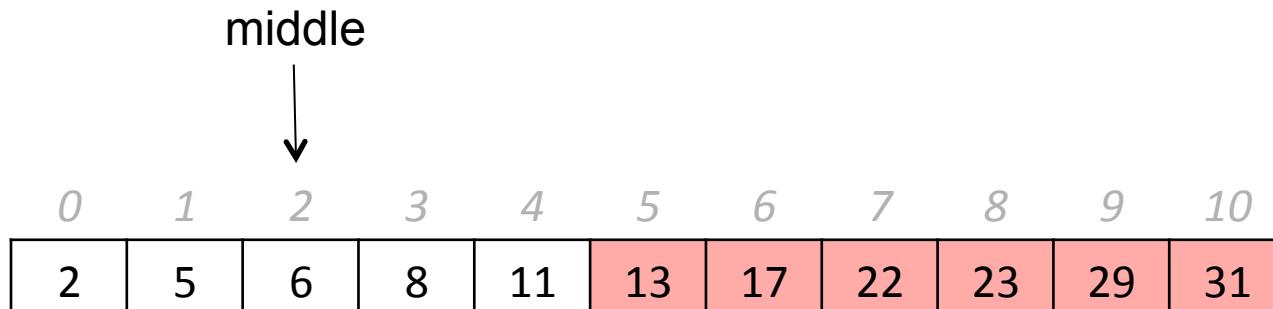
Binary Search in Action

- Search for 7:
- Look at 13
 - it's too big, so we rule out indices 5-10



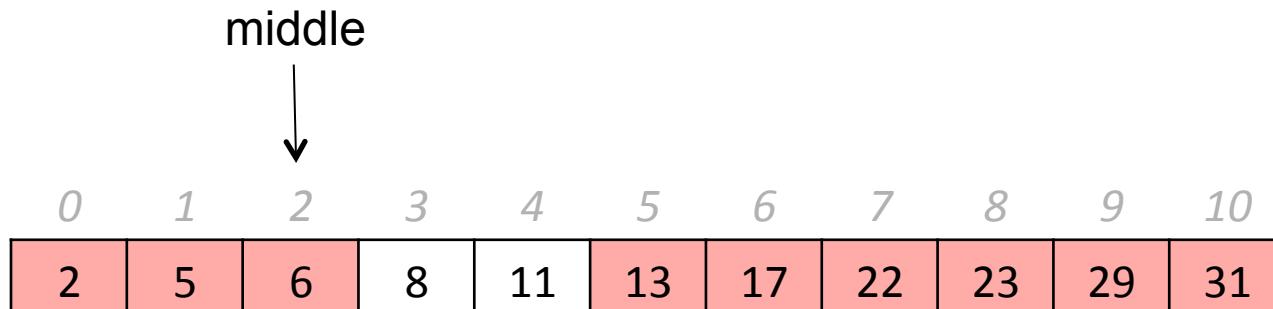
Binary Search in Action

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 - it's too big, so we rule out indices 5-10
- Pick the new middle of the remaining elements
- Look at 6:



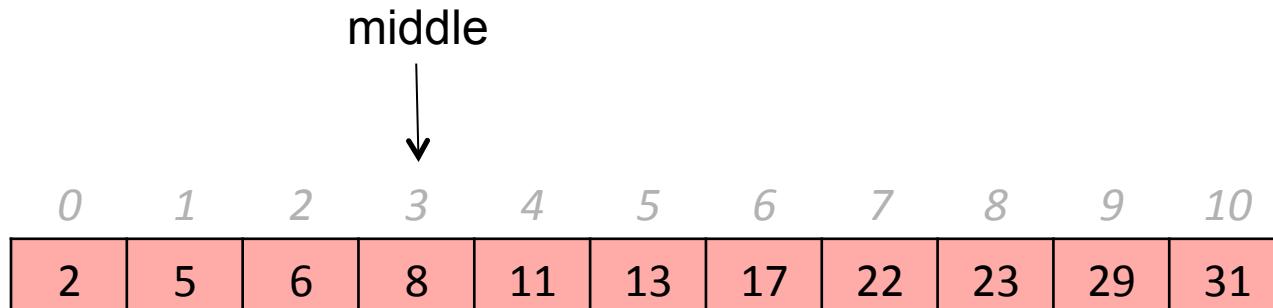
Binary Search in Action

- Search for 7:
- Look at 13
 - it's too big, so we rule out indices 5-10
- Pick the new middle of the remaining elements
- Look at 6:
 - it's too small, so we rule out indices 0-3



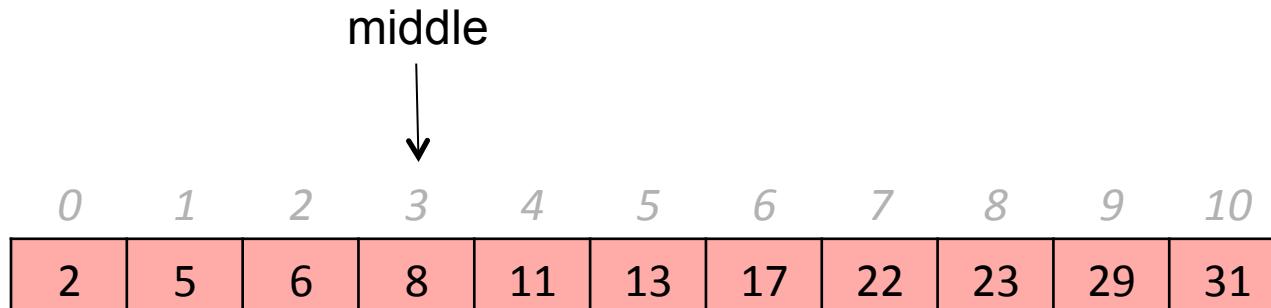
Binary Search in Action

- Search for 8:
- Look at 13
 - it's too big, so we rule out indices 5-10
- Look at 6:
 - it's too small, so we rule out indices 0-3
- Look at 8:
 - it's too big! We rule out elements 3-4



Binary Search in Action

- Search for 8:
- Look at 13
 - it's too big, so we rule out indices 5-10
- Look at 6:
 - it's too small, so we rule out indices 0-3
- Look at 8:
 - it's too big! We rule out elements 3-4
- No elements left to search – we return false



Sorted Array

- What if we **sorted** the array?
 - What is the Big-Oh of contains?
 - $O(\log N)$
 - What is the Big-Oh of adding an element?
 - $O(N)$
 - What is the Big-Oh of removing an element?
 - $O(N)$

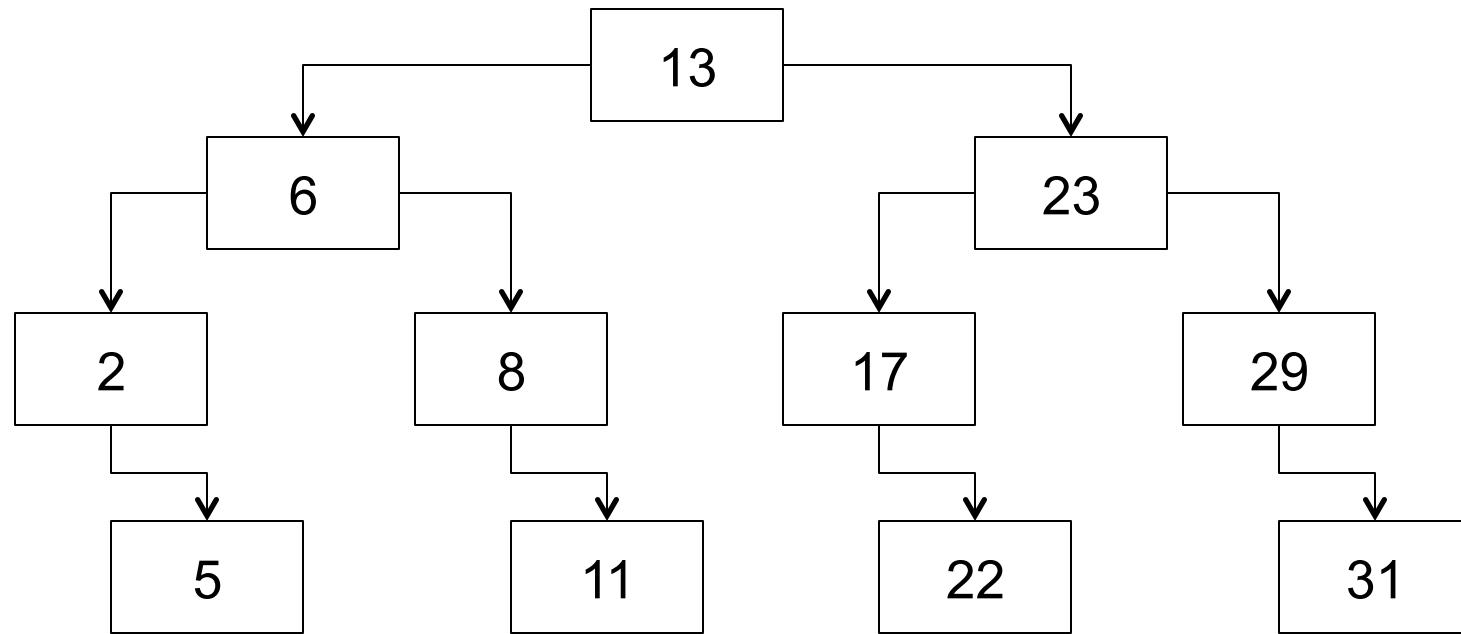
0	1	2	3	4	5	6	7	8	9	10
2	5	6	8	11	13	17	22	23	29	31

A Modification

- Problem: an array is slow to insert into or remove from
- Our solution was a **linked list** – have each element connected to one other element
 - Easy to add/remove elements
 - Can't skip elements – need to go in order
- Maybe we can find some way to implement the jumps necessary for binary search...

A Modification

- What are all the possible paths binary search could take on this array (ties are broken by choosing the smaller element)?



0	1	2	3	4	5	6	7	8	9	10
2	5	6	8	11	13	17	22	23	29	31

A Modification

- Key idea: we always jump to one of two elements in binary search (depending on if the element we're looking at is too big or too small)
- What if we had a Linked List where we stored two pointers, allowing us to make those jumps quickly?

Binary Search Tree

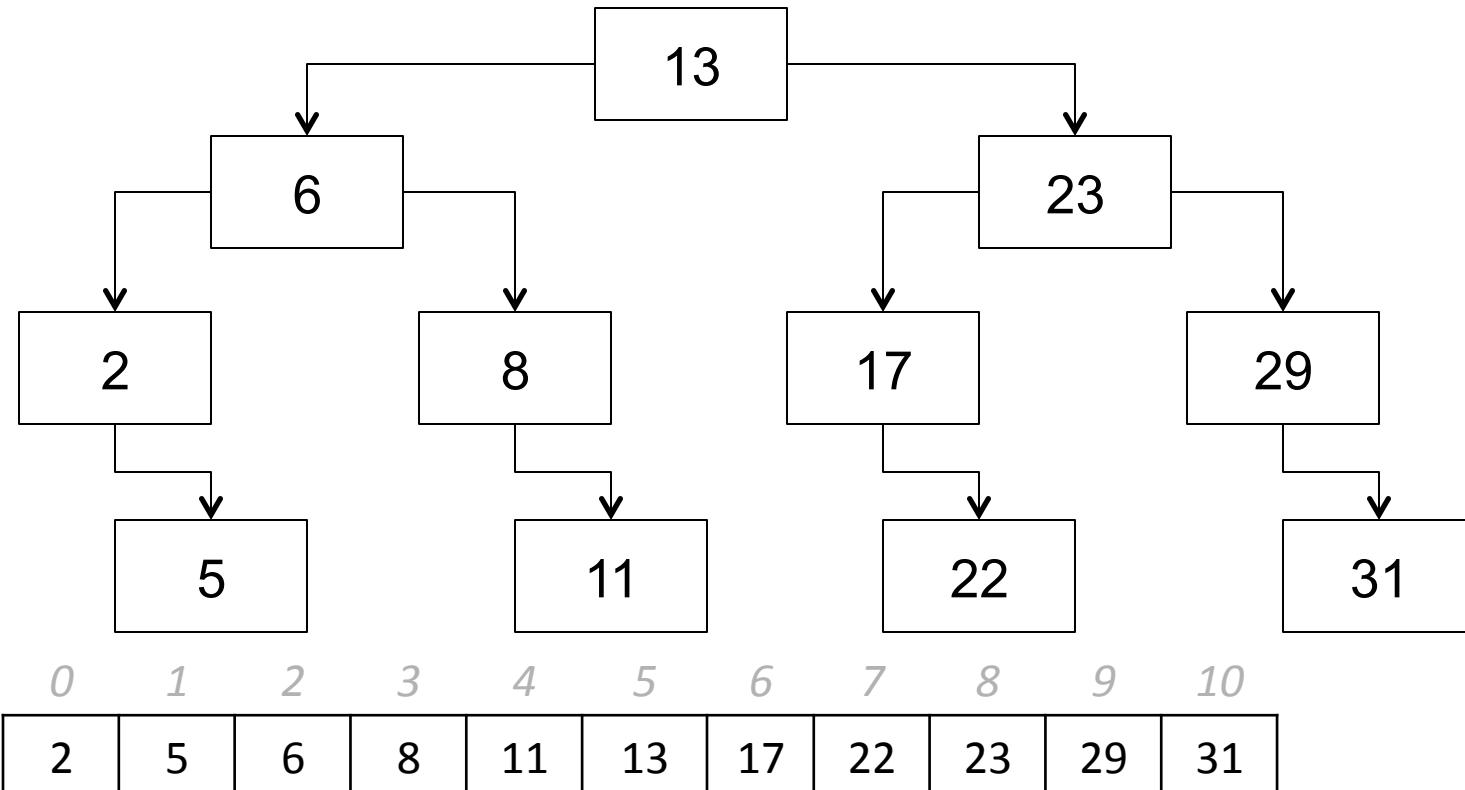
- A **tree** is a data structure where each element (**parent**) stores two or more pointers to other elements (its **children**)
 - A doubly-linked list doesn't count because, just like outside of computer science, a child can not be its own ancestor
- Each node in a **binary tree** has two pointers
 - Some of these pointers may be `nullptr` (just like in a linked list)
 - We'll see examples of non-binary trees in future lectures
- A **binary search tree** is a binary tree with special ordering properties that make it easy to do binary search
- Similar to a Linked List:
 - Each element in its own block of memory
 - Have to travel through pointers (can't skip "generations")

(Binary) TreeNode

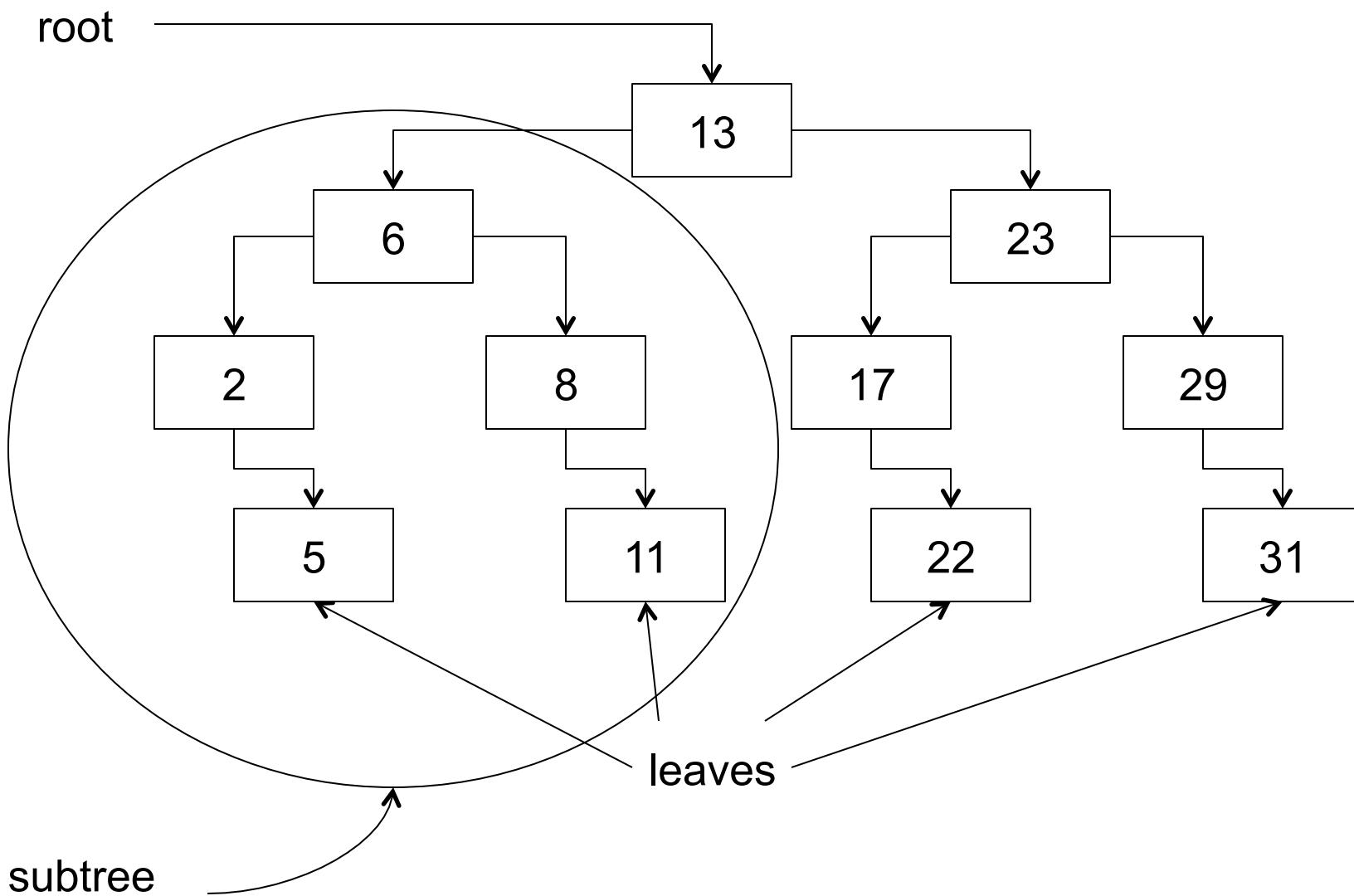
```
struct TreeNode {  
    int data; // assume that the tree stores ints  
    TreeNode *left;  
    TreeNode *right;  
};
```

Binary Search Trees

- We'll say a binary search tree has the following property:
 - All elements to the left of an element are smaller than that element
 - All elements to the right of an element are bigger than that element
 - Just like our sorted array!

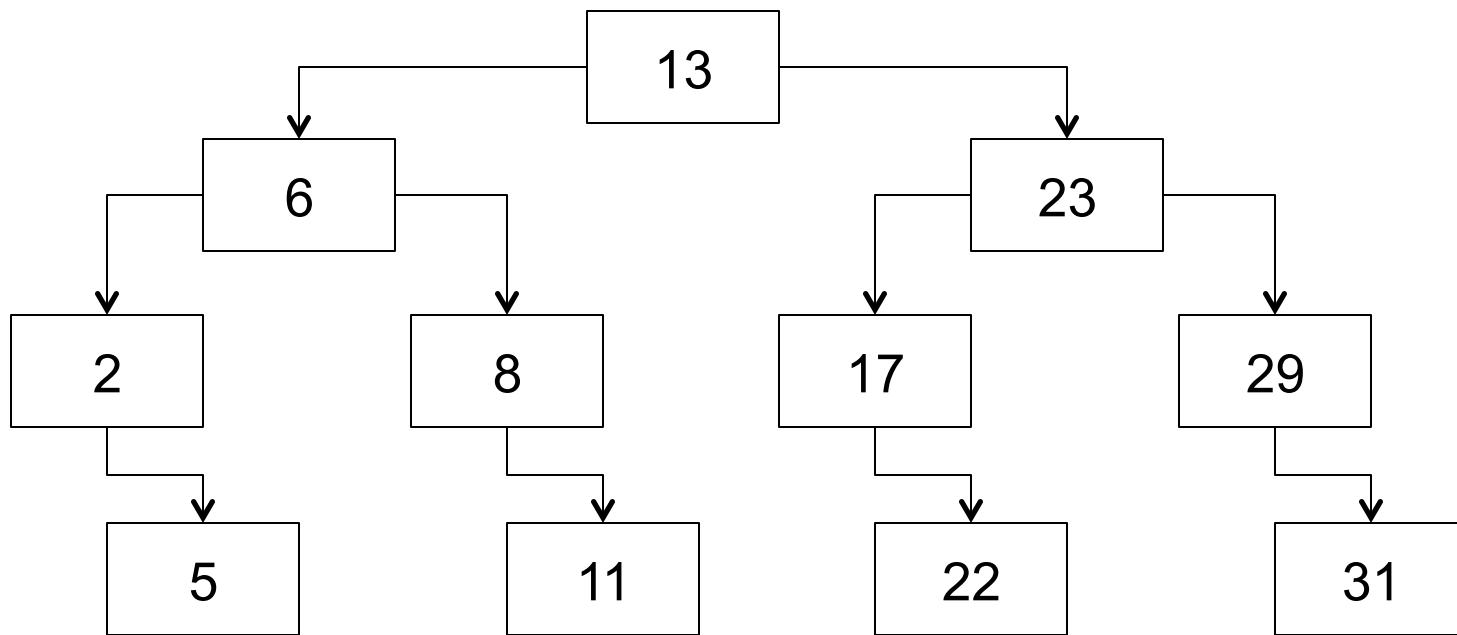


Tree anatomy



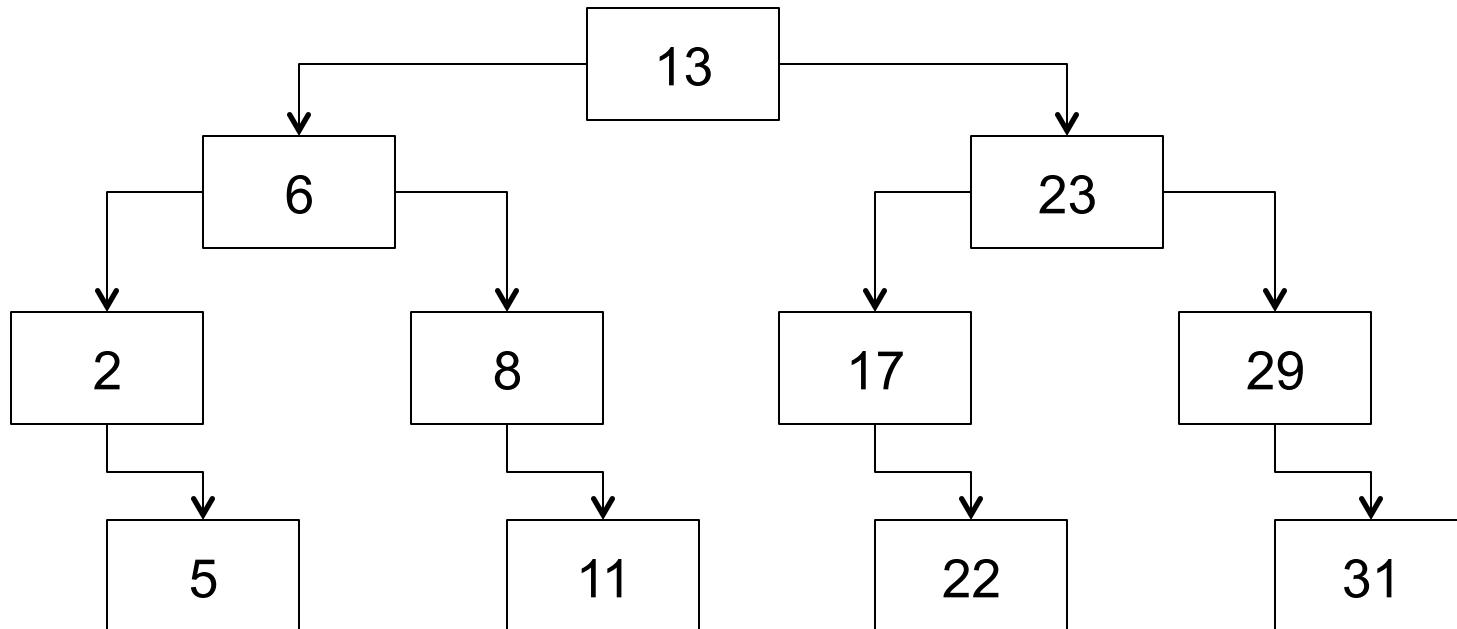
BST Contains

- How would you search a BST for an element?



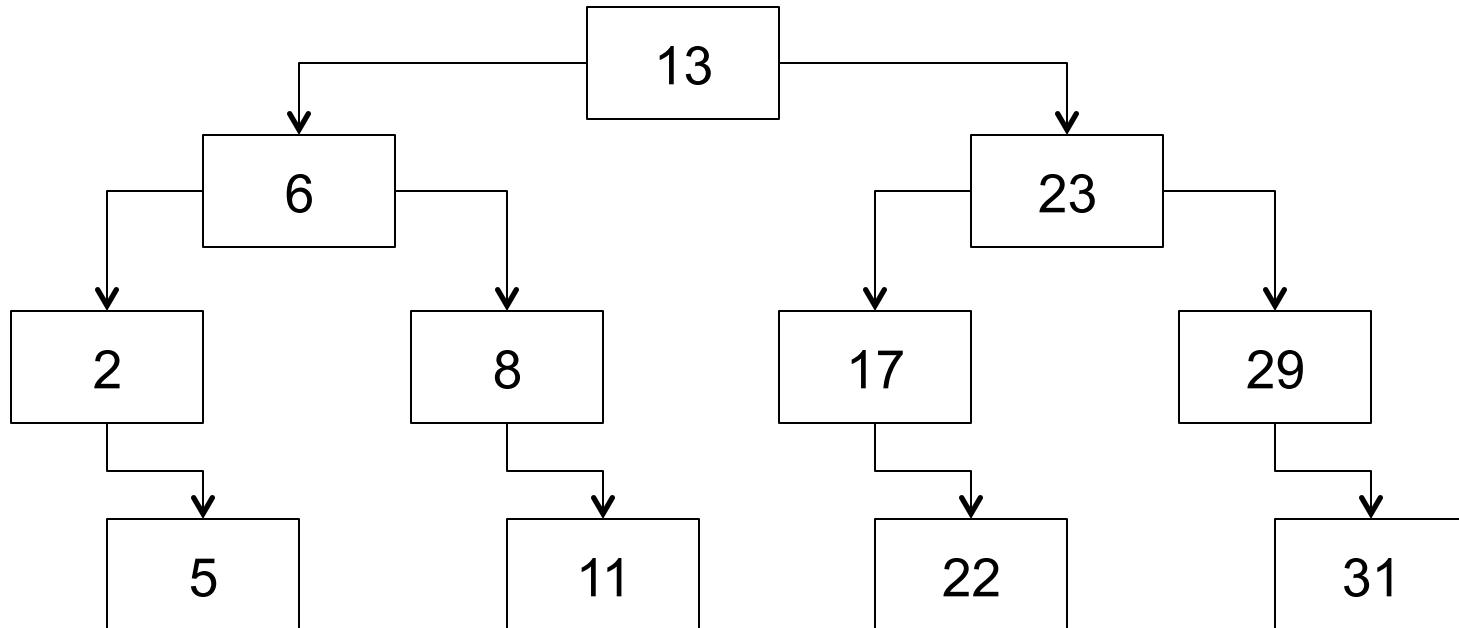
BST Contains

- How would you search a BST for an element?
- Start at root:
 - If root is too big, go left (entire right subtree is too big)
 - If root is too small, go right (entire left subtree is too small)



Trees and Recursion

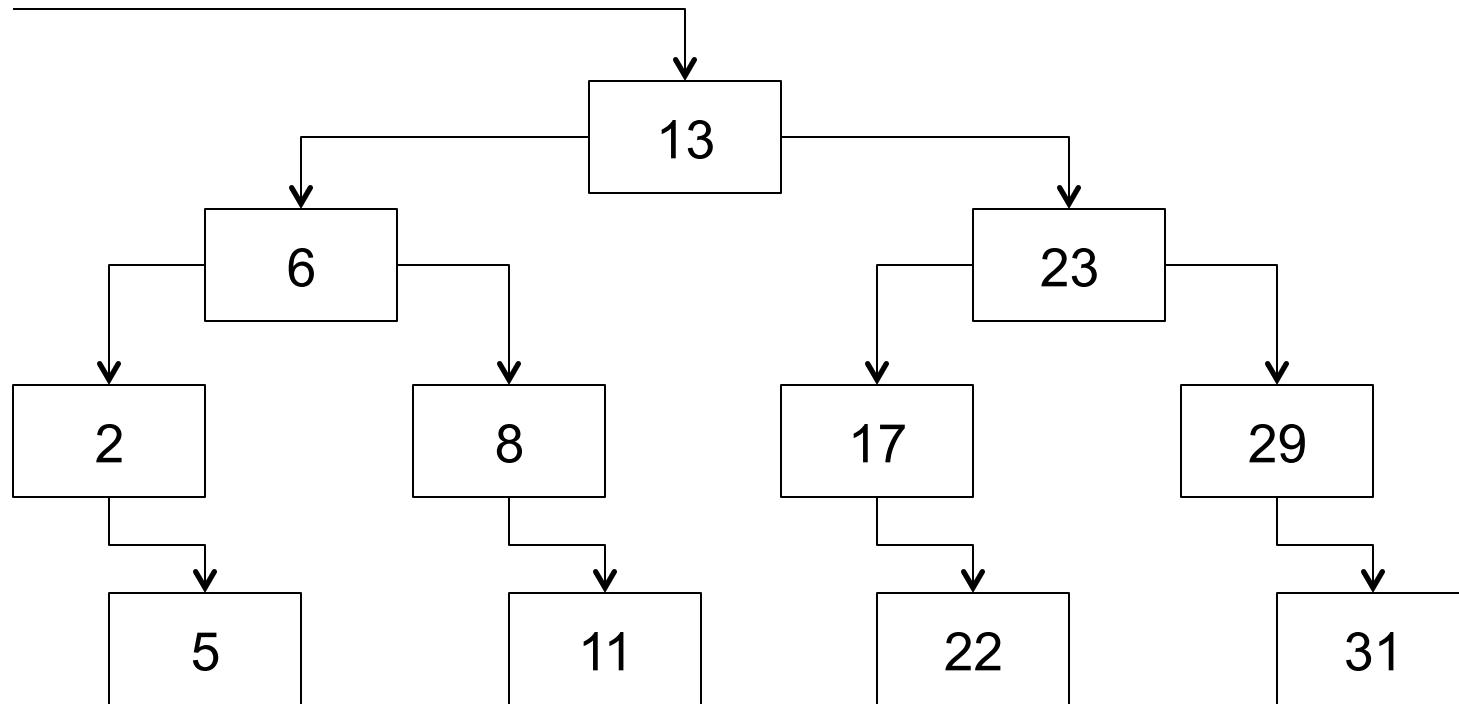
- Trees are fundamentally **recursive** (subtrees are smaller trees)
- Start at root:
 - If root is too big, go left (entire right subtree is too big)
 - If root is too small, go right (entire left subtree is too small)



Trees and Contains

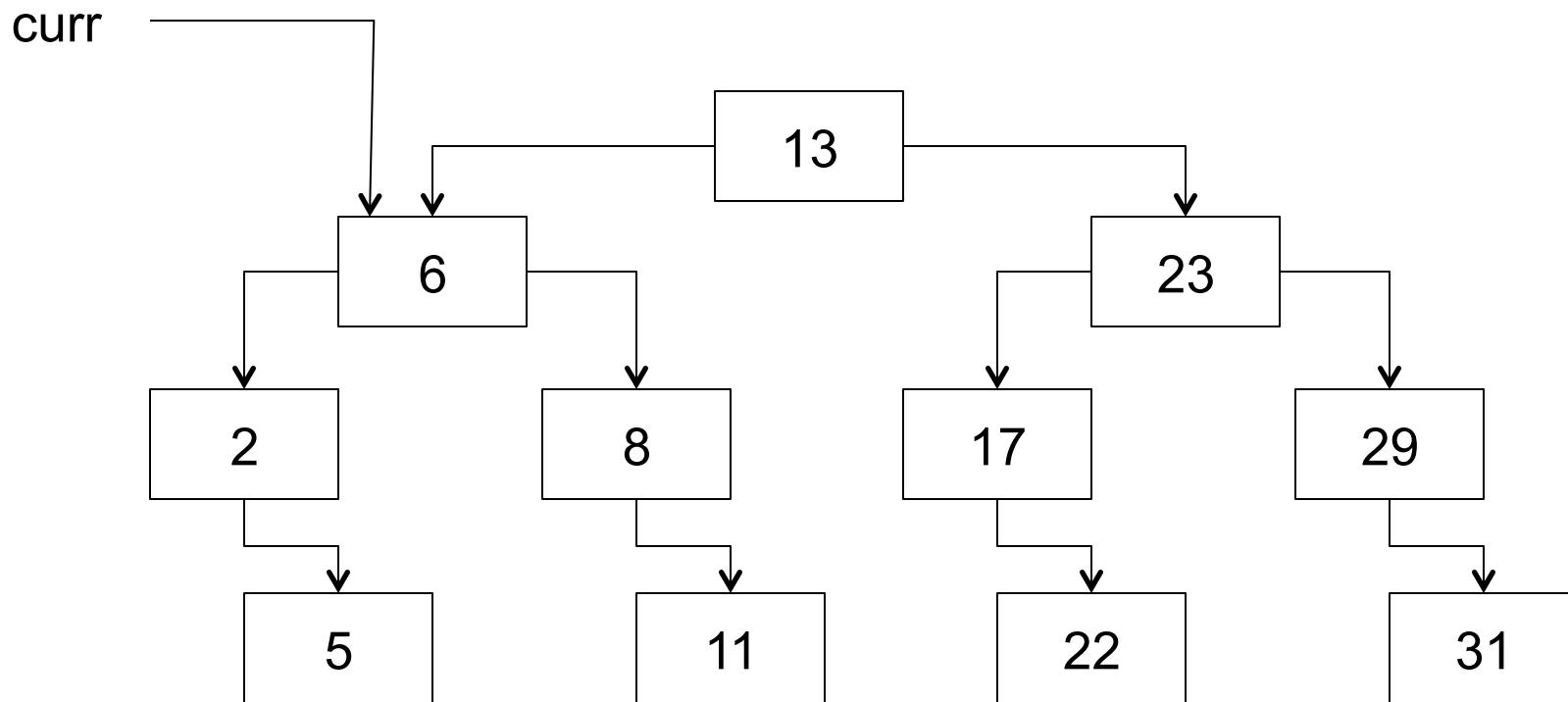
- Search for 5
- Start at root:
 - If root is too big, go left (entire right subtree is too big)
 - If root is too small, go right (entire left subtree is too small)

curr



Trees and Contains

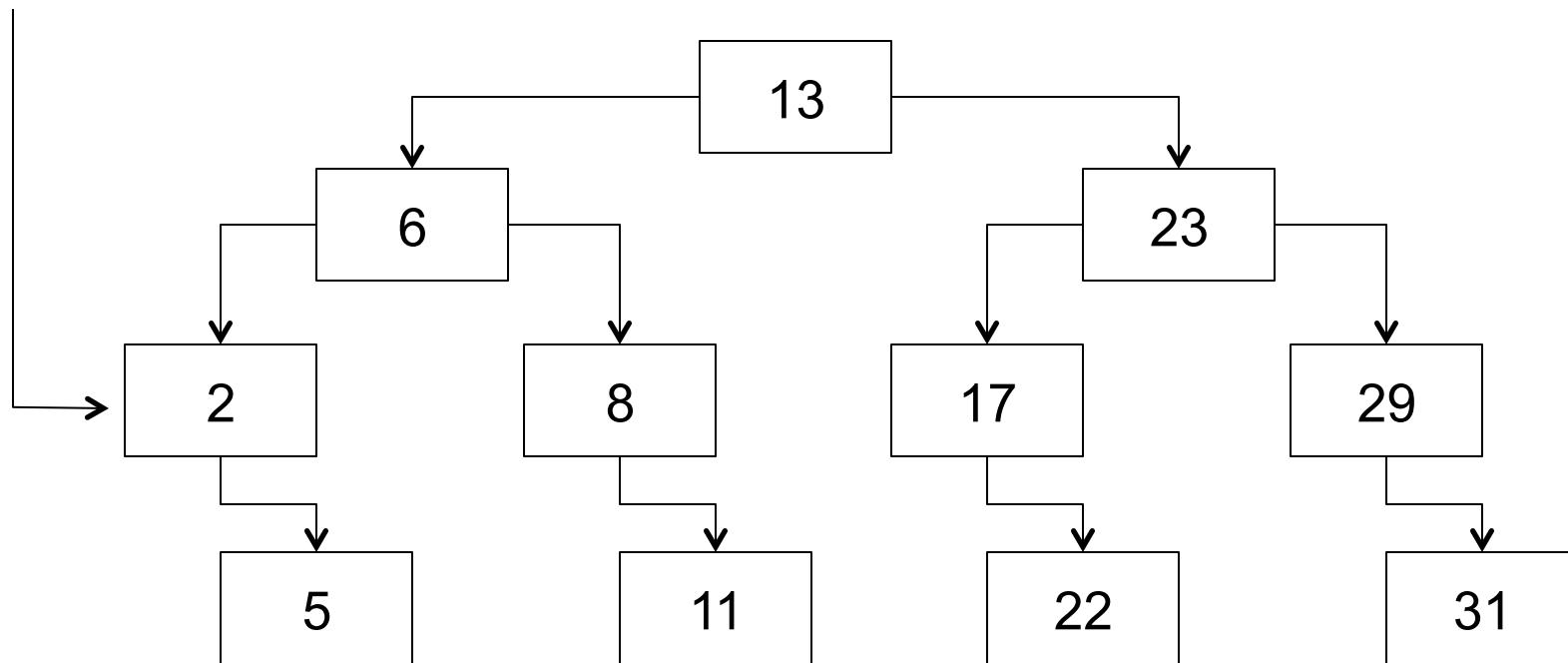
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Trees and Contains

- Search for 5
- Start at root:
 - If root is too big, go left (entire right subtree is too big)
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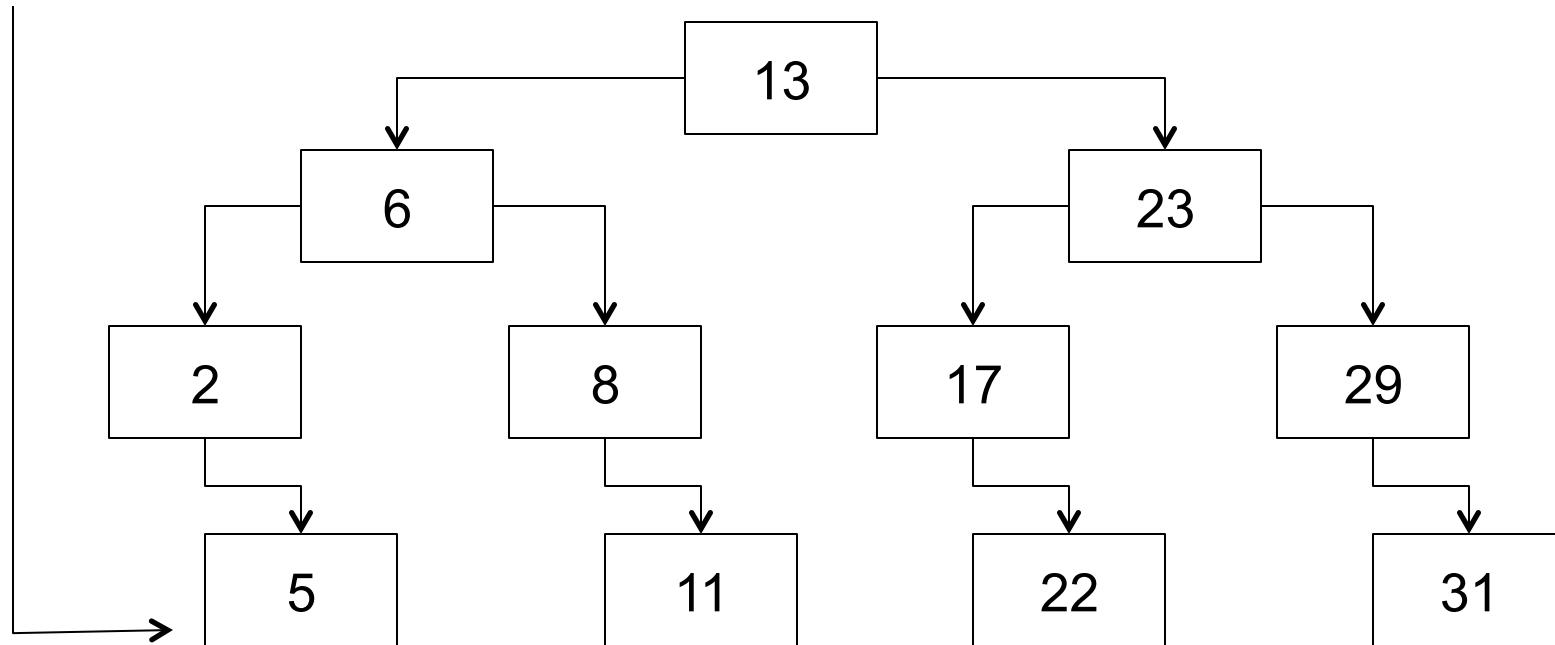
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Trees and Contains

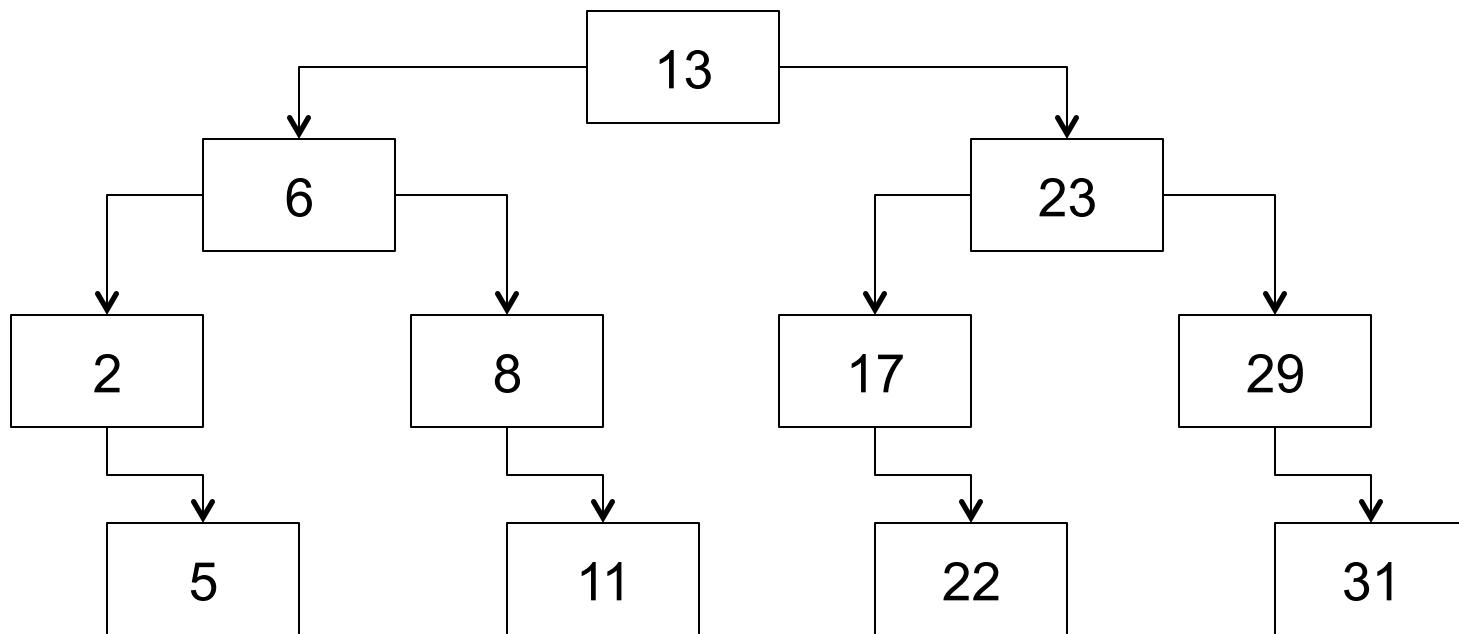
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curr



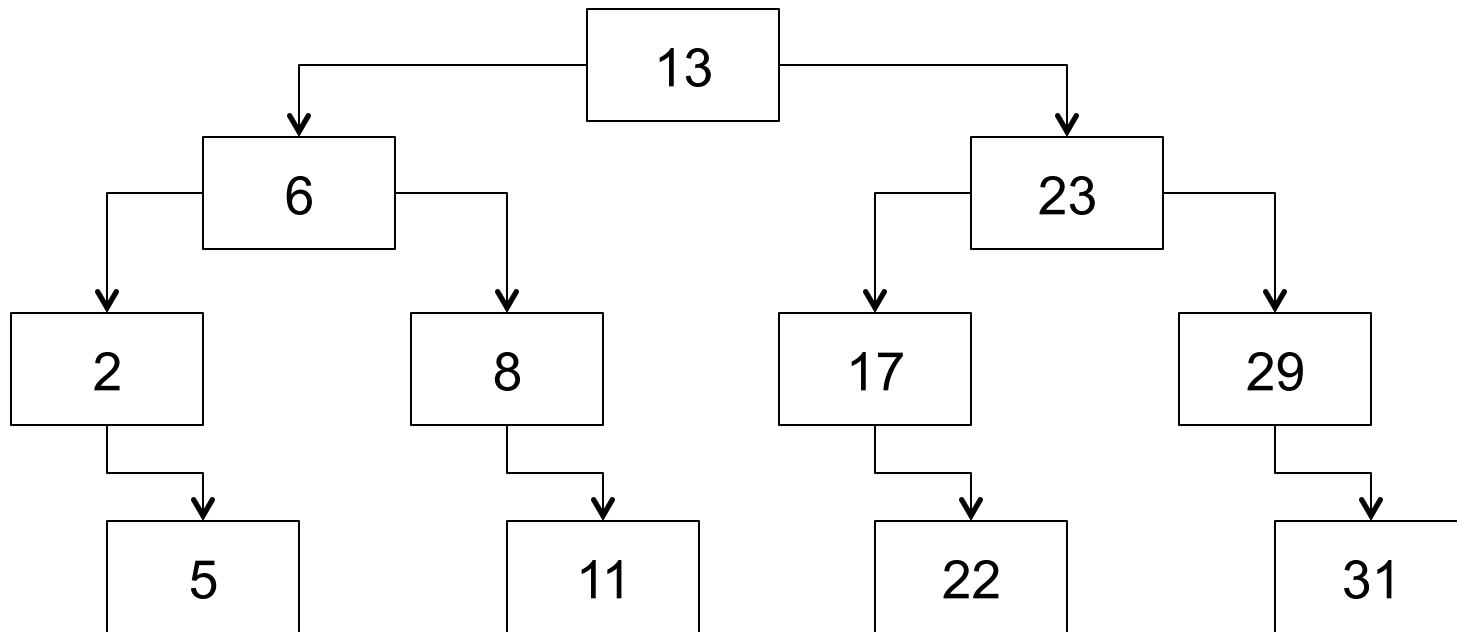
Printing Trees

- We need to be able to print our Set
- How would we print a tree?



Printing Trees

- How would we print a tree?
 - Idea: need to recurse both left and right
 - **Traverse the tree!**
 - Most tree problems involve traversing the tree

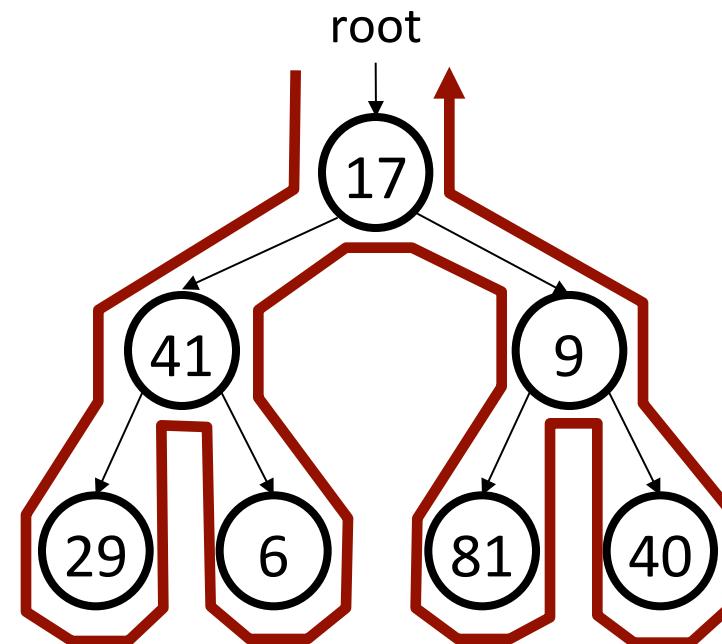


Traversal trick

- To quickly generate a traversal:
 - Trace a path counterclockwise.
 - As you pass a node on the proper side, process it.

- pre-order: left side
- in-order: bottom
- post-order: right side

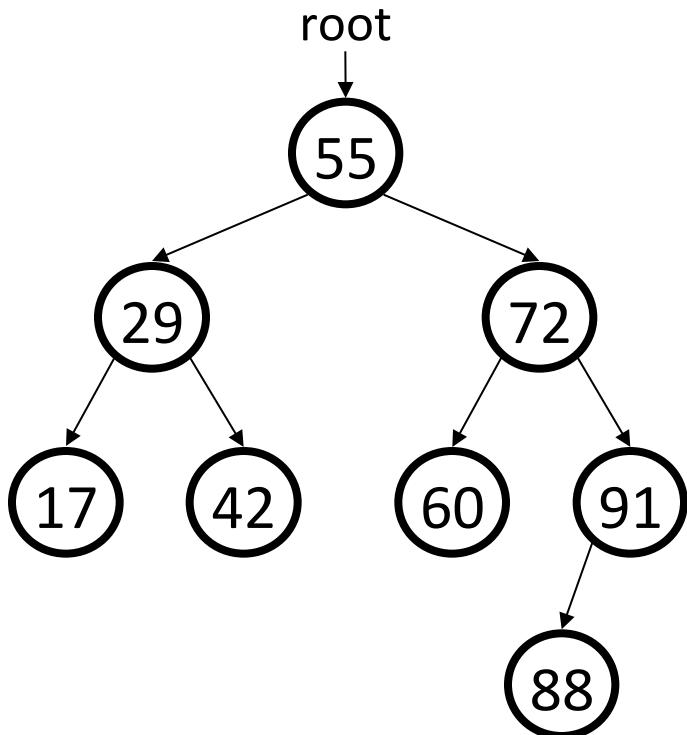
- What kind of traversal does a for-each loop in a Set do?



- pre-order: 17 41 29 6 9 81 40
- in-order: 29 41 6 17 81 9 40
- post-order: 29 6 41 81 40 9 17

getMin/getMax

- Sorted arrays can find the smallest or largest element in $O(1)$ time (how?)
- How could we get the same values in a binary search tree?



Announcements

- Assignment 4 is due **tomorrow**
- Assignment 5 will be released tomorrow
 - More time to complete it, but this assignment will be significantly longer than the others you've seen this quarter
 - As a rough guide, part c took SLs about four times as long to solve as part a, so don't wait until the last minute
- You will get assignment 3 feedback on today
- Please give feedback (if you have the next 30 minutes free):
cs198.stanford.edu
- Exam logistics
 - Midterm today, July 25, from 7:00-9:00PM in Hewlett 200

AMA

You've worked hard and have an exam today – you can leave early or stick around to ask me questions 😊