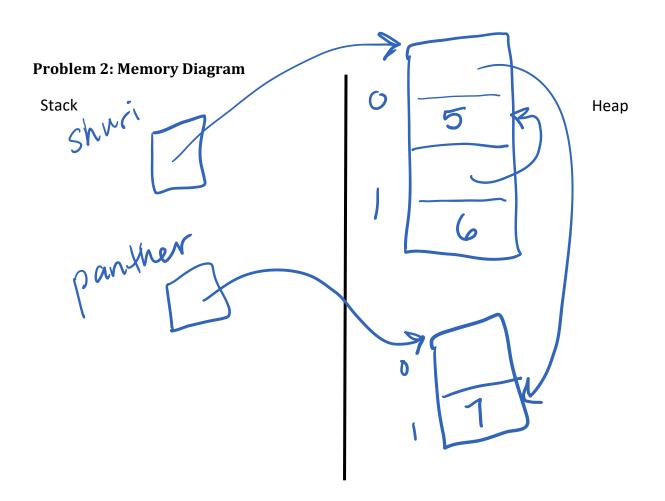
This exam is based on the final exam given in Fall 2018. The class was taught by Cynthia Lee. This was a 3-hour paper exam.

Problem 1: Floating Point

- (a) 160
- (b) 18
- (c) 176 we lose precision because minifloat does not have enough bits to store the entire sum of the number. This likely would not have been an issue with IEEE 32-bit floats, since there are many more bits to store numbers with more precision.



Problem 3: Pointers and Generics (12pts)

```
void mixup1(char *ptr1, char *ptr2) {
(a)
           swap(ptr1 + 1, ptr2 + 1, sizeof(int));
     }
     -256 = 0 \times FFFFF00
(b)
     255 = 0x000000FF // notice 0th byte of each is swapped
(c)
     //little-endian solution (this is how myth works)
(d)
     void mixup2(int *ptr1, int *ptr2) {
         swap(ptr1 + 1, ptr2 + 2, sizeof(char));
     }
     // big-endian solution (we also gave points for this even though
     //not how myth works)
     void mixup2(int *ptr1, int *ptr2) {
         swap((char*)ptr1 + 7, (char*)ptr2 + 11, sizeof(char));
     }
```

Problem 4: Assembly (23pts)

```
(a) Intended solution:
char *vp(unsigned int aaron, char *burr)
{
    unsigned int leslie = strlen(burr) * 3;
    if (aaron < -16) {
                         // or <= -17 (equivalent for unsigned)</pre>
        while (leslie > 6) {
            leslie /= 4;
        }
        if (aaron >= 256) {
            pass final level(leslie);
        } else {
            explode_bomb();
        }
    } else {
        explode bomb();
    }
    return burr + 3;
}
```

- (b) Caller-owned registers must be saved before we use them as local/temporary storage, then restored before the function returns.
- (c) The lea is faster than imul, and achieves * 3 by multiplying the value by 2 and then adding the product to itself. Shifting left by 2 with shr is the same as multiplying by 4 in binary arithmetic, and shr is faster than divide.
- (d) It may look like it is not possible, because aaron must be both less than -16 and greater than 256. However, the comparison with -16 is unsigned, because when signed and unsigned are compared the unsigned equivalents of both values are used. So we just need a value for aaron that satisfies 256 <= aaron < 0xFFFFFF0 (which is -16U). Some common solutions that would work: 256, 257, 300, -17, 0xFFFFFF0, or "max int." Some common solutions that don't work: 255, -1 (or 0xFFFFFFF or "max unsigned int"), 0xFFFFFF0.

Problem 5: Heap Allocator (24pts)

(b) Without a footer to inform us of how many bytes are in the block to the left in O(1) time (the way the header informs us of how many bytes are in the block to the right in O(1) time), we would need to traverse either the free list or the entire heap, looking for the block whose header tells us it is to our left. That looping is O(N).

```
(c)
struct Header *pay_to_hdr(struct Node *payload)
{
    // Possible solution 1
    return (struct Header *)payload - 1;
    // Possible solution 2
    return (struct Header *)((char *)payload - HDRSIZE);
}
```

```
(d)
size t count free inorder()
{
   size t nfree = 0;
   for (struct Header *curr = segment start; curr != NULL;
                       curr = get neighbor(curr)) {
      if (curr->used == 0) // need this because we look at both free and used
         nfree++;
   }
   return nfree;
}
(e)
size t count free list()
{
   size t nfree = 0;
   // needs NULL check in pay to hdr()
   for (struct Node *curr = free list; curr != NULL;
                     curr = curr->next) {
      if (pay to hdr(curr)->used == 0) // not needed because all are free
         nfree++;
   }
   return nfree;
}
     You may remove the check from count free list (part (e)), for the reasons explained in
(f)
```

the comments for (d) and (e), which is that the free_list contains only free blocks, so there is no need to check the free flag.

```
(g)
void update_header(struct Header *left)
{
    // remember to divide by 8, because stored as number of 8-byte words
    left->nwords += get_neighbor(left)->nwords + HDRSIZE / 8;
}
```

```
(h)
void remove_node(struct Node *remove)
{
    if (remove->prev == NULL) {
       free_list = remove->next;
    } else {
       remove->prev->next = remove->next;
    }
    if (remove->next == NULL) {
       // Empty because no later node needs its prev set.
       ;
    } else {
       remove->next->prev = remove->prev;
    }
}
```