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# Lecture 10:

## Skew Tolerant Domino Clocking

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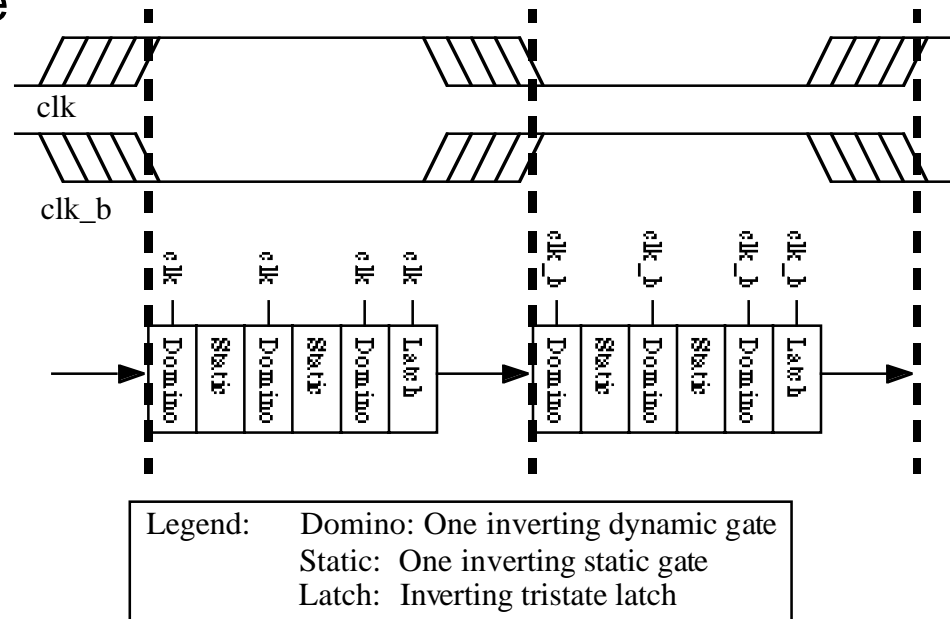
# Introduction

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- Domino Circuits are becoming ubiquitous in high speed digital ICs
  - Offer 30% (or more) speedup over static CMOS raw gate delay
  - Dual-rail domino becoming more common because many functions are nonmonotonic, area is less of an issue
- Nevertheless, traditional domino pipelines have significant overhead
  - Latch required to hold result while next stage evals, prev. precharges
  - Skew budget, no time borrowing, latch delay
- Look at several ways to reduce this overhead
  - Better latches, Self-timing
  - Skew-tolerant domino is a powerful new technique
- Evaluate performance benefits of skew-tolerant domino

# Domino from a System Perspective

- Domino doesn't look so attractive in the context of a traditional pipeline



1. Pay clock skew twice each phase
2. Balancing short phases is hard since there is no time borrowing
3. Latches become a significant fraction of the cycle time

# Traditional Domino Performance Evaluation

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- Let  $T = \text{cycle time} = 20 \text{ FO4 delays}$ ;  $t_{\text{skew}} = 2$ ;  $t_{\text{setup}} = 1^1$ 
  - Difficult filling cycle exactly (no time borrowing)  $\rightarrow t_{\text{imbalance}} = 1$
- $T_{\text{phase-logic}} = T/2 - t_{\text{skew}} - t_{\text{setup}} - t_{\text{imbalance}}$
- Baseline Design:
  - $T_{\text{phase-logic}} = 20/2 - 2 - 1 - 1 = 6$
  - 40% of the phase is wasted in overhead! Slower than static!
- Optimized Design:
  - Define clock domains and use  $t_{\text{skew-local}} = 1$
  - Work hard to balance logic between phases:  $t_{\text{imbalance}} = 0$  (optimistic)
  - $T_{\text{phase-logic}} = 20/2 - 1 - 1 - 0 = 8$
  - Still, 20% of the phase is overhead!

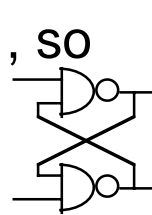
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1. Remember for this situation, the setup time must be large enough that the output has settled before clock arrives since the output might go into a dynamics gate on the next cycle and might not be monotonic

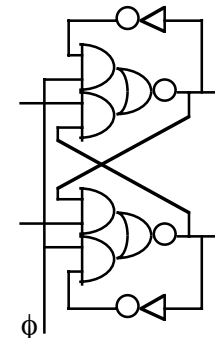
# Early Enhancements

- Good designers have recognized this problem for years.
- The largest problem is the hard edges set by the latches.
- A variety of latches soften this edge:

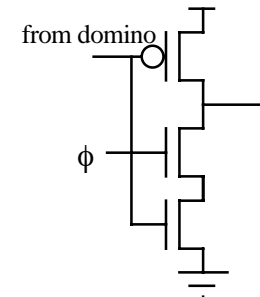
- Gate outputs are already  $\_q1$ , so why use another clock. An SR latch will work instead
- Use the monotonic nature of the signal to feed it into a precharged latch stage



SR Latch



Dual-Monotonic Latch



TSPC Latch

- Still have a problem if you want to use non-monotonic logic somewhere, since logic must settle before earliest clock, while gate might not evaluate until a late clock
- But if you only have monotonic gates ...

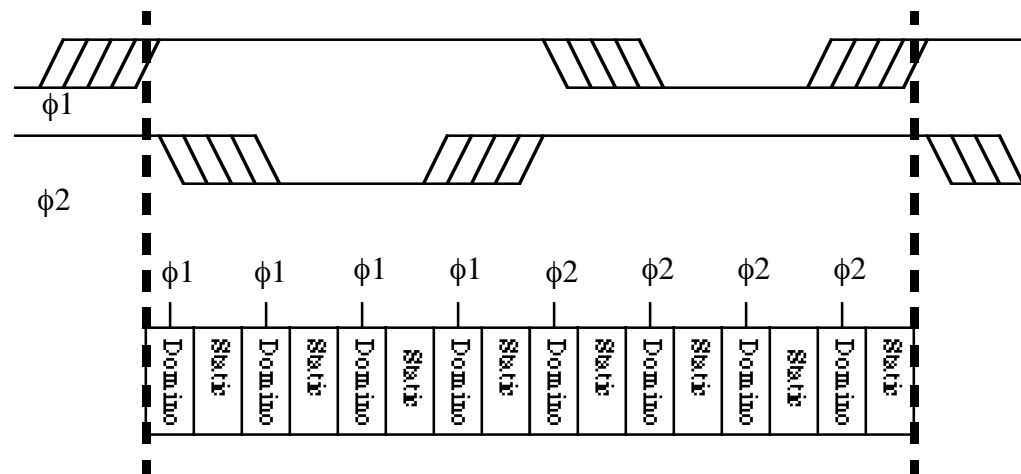
# Skew Tolerant Domino Clocking

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- If inputs are all dual rail, then as long as the clock arrives before the data,
    - The gate will wait and fire when the data arrives
  - If the next gate fires before the current gate precharges,
    - There is no need for a latch
  - Like the self-timed pipeline
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- Can generate these properties using overlapping clocks

# Skew-Tolerant Domino Circuits

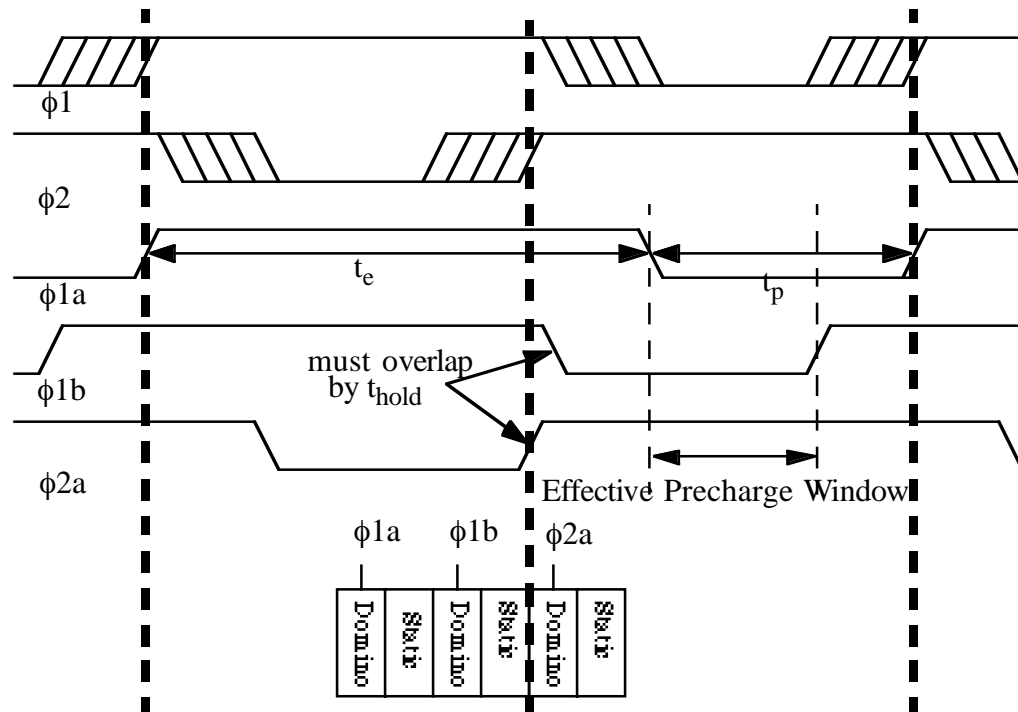
- How much clock skew could we tolerate given N clock phases?
  - Divide logic into N phases of  $T/N$  duration each.
  - Overlapping clocks eliminates need for latches
  - Extra overlap accommodates clock skew and time borrowing



- As with other domino techniques, budget skew on the transition from static to domino

# Skew Tolerance

- $T = t_e + t_p$
- $t_p = t_{prech} + t_{skew}$ ;  $t_e = T/N + t_{skew} + t_{hold}$
- Hence  $t_{skew-max} = [T(N-1)/N - t_{prech} - t_{hold}] / 2$



# Numerical Example

- Let  $t_{\text{prech}} = 4$ , long enough to:
- precharge domino gate
- make subsequent skewed static fall below  $V_t$
- $t_{\text{hold}}$  is slightly negative for reasonable cell libraries
- next phase can evaluate before precharge ripples through static gate
- conservatively bound  $t_{\text{hold}}$  at 0

N	$t_{\text{skew}}$	$t_p$
2	2	6
3	3.33	7.33
4	4	8
6	4.66	8.66
8	5	9

- Sweet spots: N=2 (fewest clocks), N=4 (good tolerance, 50% duty cycle)

# Global & Local Skew

- This is good, but we can do better!
- Local skew can be more tightly controlled than global skew (~ 1 FO4)
  - Require that each phase of logic fit in a local clock domain:
- $t_p = t_{\text{prech}} + t_{\text{skew-local}}$ ;  $t_e = T/N + t_{\text{skew-global}} + t_{\text{hold}}$
- Hence  $t_{\text{skew-global-max}} = T(N-1)/N - t_{\text{skew-local}} - t_{\text{prech}} - t_{\text{hold}}$
- When  $t_{\text{skew-global}}$  gets huge, precharge interferes with subsequent phase

N	$t_{\text{skew-global}}$	$t_p$
2	2	5
3	5.66	5
4	6	6
6	6	7.33
8	6	8

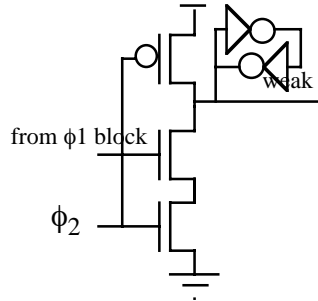
# Time Borrowing

- We don't need such a large global skew tolerance!
  - Use some of this time instead to allow time borrowing
- $t_{\text{borrow}} = T(N-1)/N - t_{\text{skew-global}} - t_{\text{skew-local}} - t_{\text{prech}} - t_{\text{hold}}$
- Intentional borrowing helps balance logic between phases
- Opportunistic time borrowing compensates for uncertainties in models, analysis tools, and processing
- If actual  $t_{\text{skew-global}} = 2$ ,  $t_{\text{skew-local}} = 1$ :

N	$t_{\text{borrow}}$	$t_p$
2	1	5
3	3.66	5
4	5	5
6	6.33	5
8	7	5

# Other Design Issues

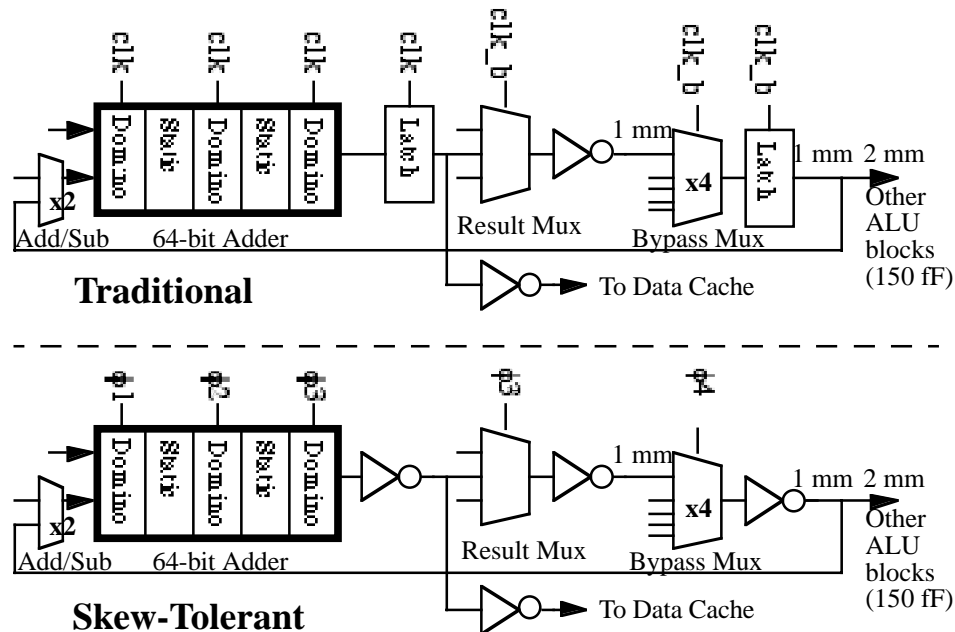
- State is no longer stored in the latch at the end of a phase
  - Instead, it is held by the first domino gate in the phase
  - Use a “full keeper” to allow stop-clock operation



- All systems with overlapping clocks require min-delay checks
  - Domino paths are presumably critical anyway, so few min-delay errors
  - 4-phase has effectively no min-delay risk
    - Overlap of all four phases is at most very small
    - A minimum of 8 gates are in the cycle anyway

# Skew-Tolerant Performance Evaluation

- Evaluate ALU self-bypass of superscalar  $\mu$ proc (like DEC Alpha)
- 3-metal 0.6  $\mu$ m process
- FO4 delay in TT corner = 138 ps
- Compare traditional domino to 4-phase skew-tolerant domino



# Simulation Results

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- No Skew:
  - Traditional Domino: Latency = 13.0 FO4, cycle time = 16.6
    - Cycles are unbalanced; no time borrowing available
  - Skew-Tolerant Domino: Latency = 11.9 FO4, cycle time = 11.9
    - Remove latches from critical path, balance pipe stages
- 1 FO4 local skew:
  - Traditional Domino: Latency = 15.0 FO4, cycle time = 17.6
    - Skew adds to both phases for latency
    - Unbalanced second stage already has margin in the cycle time
  - Skew-Tolerant Domino: Latency = 11.9 FO4, cycle time = 11.9
    - Skew is tolerated

# Summary

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- Offers most of the benefits of self-timed designs while preserving the simplicity of a synchronous methodology.
- Clock generation & distribution becomes key issue. However, control generation and distribution can be just as tough in self-timed designs.
- Skew-Tolerant Domino eliminates most of the overhead found in traditional domino systems:
  - Tolerates clock skew
  - Removes latches from the critical path
  - Allows time borrowing
  - Robust
- High-performance microprocessor designs have used these ideas but they don't talk about them.