

Mathematical Logic

Part One

Question: How do we formalize the definitions and reasoning we use in our proofs?

Where We're Going

- ***Propositional Logic*** (Today)
 - Basic logical connectives.
 - Truth tables.
 - Logical equivalences.
- ***First-Order Logic*** (Wednesday/Friday)
 - Reasoning about properties of multiple objects.

Propositional Logic

A ***proposition*** is a statement that is,
by itself, either true or false.

Some Sample Propositions

- Puppies are cuter than kittens.
- Kittens are cuter than puppies.
- Usain Bolt can outrun everyone in this room.
- CS103 is useful for cocktail parties.
- This is the last entry on this list.

Propositional Logic

- ***Propositional logic*** is a mathematical system for reasoning about propositions and how they relate to one another.
- Every statement in propositional logic consists of ***propositional variables*** combined via ***propositional connectives***.
 - Each variable represents some proposition, so each variable has value true or false.
 - Connectives encode how propositions are related.

Propositional Connectives

- **Logical NOT: $\neg p$**
 - $\neg p$ is true if and only if p is false.
 - Also called *logical negation*.
- **Logical AND: $p \wedge q$**
 - $p \wedge q$ is true if and only if both p and q are true.
 - Also called *logical conjunction*.
- **Logical OR: $p \vee q$**
 - $p \vee q$ is true if and only if at least one of p or q are true (inclusive OR)
 - Also called *logical disjunction*.

Truth Tables

- A ***truth table*** is a table showing the truth value of a propositional logic formula as a function of its inputs.
- Useful for several reasons:
 - They give a formal definition of what a connective “means.”
 - They give us a mechanical way to evaluate a complex propositional formula.

Truth Table for XOR

- Recall that our OR connective is inclusive.
- The truth table at right defines an exclusive or called XOR.
- We also could have expressed XOR using just the connectives we already had.

p	q	p XOR q
F	F	F
F	T	T
T	F	T
T	T	F

Which expresses XOR?

- (A) $(p \wedge q) \vee (p \vee q)$
- (B) $(p \wedge q) \vee \neg(p \vee q)$
- (C) $(p \vee q) \wedge \neg(p \wedge q)$
- (D) $(p \wedge q) \wedge (p \vee q)$

Answer at [PollEv.com/cs103](https://www.pollEv.com/cs103) or
text **CS103** to **22333** once to join, then **A**, **B**, or **C**.

Mathematical Implication

Truth Table for $p \rightarrow q$ (implies)

What is the correct truth table for implication? Enter your guess as a list of four values to fill in the rightmost column of the table.
(ex: F, T, ?, F)

p	q	$p \rightarrow q$
F	F	
F	T	
T	F	
T	T	

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Truth Table for Implication

p	q	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T

Truth Table for Implication

p	q	$p \rightarrow q$	
F	F	T	• Bad bracket, didn't get A
F	T	T	• Bad bracket, got A
T	F	F	• Perfect bracket, didn't get A
T	T	T	• Perfect bracket, got A

Truth Table for Implication

p	q	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T

The implication is only false if p is true and q isn't. It's true otherwise.

Truth Table for Implication

p	q	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T

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You will need to commit this table to memory.
(Consider a tattoo on your forearm.)
We're going to be using it *a lot* over the rest of the week.

Why This Truth Table?

- The truth values of the \rightarrow are the way they are because they're *defined* that way.
- Are there other ways we could write a proposition that has the same truth table as \rightarrow , using the other connectives? (*like what we did for XOR*)
 - Yep!
 - Try to think of some
 - What's the truth table for $\neg(p \wedge \neg q)$?

The Biconditional Connective

The Biconditional Connective

- The biconditional connective \leftrightarrow is used to represent a two-directional implication.
- Specifically, $p \leftrightarrow q$ means both that $p \rightarrow q$ and that $q \rightarrow p$.
- Based on that, what should its truth table look like?
- Take a guess, and talk it over with your neighbor!

Biconditionals

- The ***biconditional*** connective $p \leftrightarrow q$ is read “ p if and only if q .”
- Here's its truth table:

p	q	$p \leftrightarrow q$
F	F	T
F	T	F
T	F	F
T	T	T

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F	T	F
T	F	F
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One interpretation of \leftrightarrow is to think of it as equality: the two propositions must have equal truth values.

True and False

- In addition to variables and connectives, we have constants: true and false.
 - The symbol \top is a value that is always true.
 - The symbol \perp is value that is always false.
- These are often called connectives, though they don't connect anything.
 - (Or rather, they connect zero things.)

Proof by Contradiction

- Suppose you want to prove p is true using a proof by contradiction.
- The setup looks like this:
 - Assume p is false.
 - Derive something that we know is false.
 - Conclude that p is true.
- In propositional logic:

$$(\neg p \rightarrow \perp) \rightarrow p$$

Operator Precedence

- How do we parse this statement?

$$\neg x \rightarrow y \vee z \rightarrow x \vee y \wedge z$$

- Operator precedence for propositional logic:

\neg

\wedge

\vee

\rightarrow

\leftrightarrow

- All operators are right-associative.
- We can use parentheses to disambiguate.

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∨

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$$(\neg x) \rightarrow y \vee z \rightarrow x \vee (y \wedge z)$$

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Operator Precedence

- The main points to remember:
 - \neg binds to whatever immediately follows it.
 - \wedge and \vee bind more tightly than \rightarrow .
- We will commonly write expressions like $p \wedge q \rightarrow r$ without adding parentheses.
- For more complex expressions, we'll try to add parentheses.
- Confused? Just ask!

The Big Table

Connective	Read As	C++ Version	Fancy Name
\neg	“not”	!	Negation
\wedge	“and”	&&	Conjunction
\vee	“or”		Disjunction
\rightarrow	“implies”	<i>see PS2!</i>	Implication
\leftrightarrow	“if and only if”	<i>see PS2!</i>	Biconditional
\top	“true”	true	Truth
\perp	“false”	false	Falsity

Recap So Far

- A ***propositional variable*** is a variable that is either true or false.
- The ***propositional connectives*** are
 - Negation: $\neg p$
 - Conjunction: $p \wedge q$
 - Disjunction: $p \vee q$
 - Implication: $p \rightarrow q$
 - Biconditional: $p \leftrightarrow q$
 - True: \top
 - False: \perp

Translating into Propositional Logic

Some Sample Propositions

a: I will be in the path of totality.

b: I will see a total solar eclipse.

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“I won't see a total solar eclipse
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$$\neg a \rightarrow \neg b$$

“ p if q ”

translates to

$$q \rightarrow p$$

It does *not* translate to

$$p \rightarrow q$$

“ p , but q ”

translates to

$p \wedge q$

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c: There is a total solar eclipse today.

Some Sample Propositions

a: I will be in the path of totality.

b: I will see a total solar eclipse.

c: There is a total solar eclipse today.

“If I will be in the path of totality, but there's no solar eclipse today, I won't see a total solar eclipse.”

Which is equivalent to the sentence at left?

- (A) $a \wedge \neg c \wedge \neg b$
- (B) $(a \wedge \neg c) \rightarrow \neg b$
- (C) $a \rightarrow (\neg c \wedge \neg b)$
- (D) $(\neg a \vee c) \vee \neg b$

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a: I will be in the path of totality.

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“If I will be in the path of totality, but there's no solar eclipse today, I won't see a total solar eclipse.”

$$a \wedge \neg c \rightarrow \neg b$$

Propositional Equivalences

Quick Question:

What would I have to show you to convince you that the statement $p \wedge q$ is false?

Quick Question:

What would I have to show you to convince you that the statement $p \vee q$ is false?

De Morgan's Laws

- Using truth tables, we concluded that

$$\neg(p \wedge q)$$

is equivalent to

$$\neg p \vee \neg q$$

- We also saw that

$$\neg(p \vee q)$$

is equivalent to

$$\neg p \wedge \neg q$$

- These two equivalences are called ***De Morgan's Laws***.

De Morgan's Laws in Code

- **Pro tip:** Don't write this:

```
    if (! (p() && q())) {  
        /* ... */  
    }
```

- Write this instead:

```
    if (!p() || !q()) {  
        /* ... */  
    }
```

- (This even short-circuits correctly!)

Logical Equivalence

- Because $\neg(p \wedge q)$ and $\neg p \vee \neg q$ have the same truth tables, we say that they're **equivalent** to one another.
- We denote this by writing

$$\neg(p \wedge q) \equiv \neg p \vee \neg q$$

- The \equiv symbol is not a connective.
 - The statement $\neg(p \wedge q) \leftrightarrow (\neg p \vee \neg q)$ is a propositional formula. If you plug in different values of p and q , it will evaluate to a truth value. It just happens to evaluate to true every time.
 - The statement $\neg(p \wedge q) \equiv \neg p \vee \neg q$ means “these two formulas have exactly the same truth table.”
- In other words, the notation $\varphi \equiv \psi$ means “ φ and ψ always have the same truth values, regardless of how the variables are assigned.”

An Important Equivalence

- Earlier, we talked about the truth table for $p \rightarrow q$. We chose it so that

$$p \rightarrow q \equiv \neg(p \wedge \neg q)$$

- Later on, this equivalence will be incredibly useful:

$$\neg(p \rightarrow q) \equiv p \wedge \neg q$$

Another Important Equivalence

- Here's a useful equivalence. Start with

$$p \rightarrow q \equiv \neg(p \wedge \neg q)$$

- By De Morgan's laws:

$$p \rightarrow q \equiv \neg(p \wedge \neg q)$$

$$\equiv \neg p \vee \neg\neg q$$

$$\equiv \neg p \vee q$$

- Thus $p \rightarrow q \equiv \neg p \vee q$

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$$\equiv \neg p \vee \neg \neg q$$

$$\equiv \neg p \vee q$$

- Thus $p \rightarrow q \equiv \neg p \vee q$

If p is false, then $\neg p \vee q$ is true. If p is true, then q has to be true for the whole expression to be true.

One Last Equivalence

The Contrapositive

- The contrapositive of the statement

$$p \rightarrow q$$

is the statement

$$\neg q \rightarrow \neg p$$

- These are logically equivalent, which is why proof by contrapositive works:

$$p \rightarrow q \quad \equiv \quad \neg q \rightarrow \neg p$$

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$$x < 8 \wedge y < 8 \rightarrow x + y \neq 16$$

“If $x < 8$ and $y < 8$, then $x + y \neq 16$ ”

Theorem: If $x + y = 16$, then $x \geq 8$ or $y \geq 8$.

Proof: By contrapositive. We will prove that if $x < 8$ and $y < 8$, then $x + y \neq 16$. Let x and y be arbitrary numbers such that $x < 8$ and $y < 8$.

Note that

$$\begin{aligned}x + y &< 8 + y \\ &< 8 + 8 \\ &= 16.\end{aligned}$$

This means that $x + y < 16$, so $x + y \neq 16$, which is what we needed to show. ■

Why This Matters

- Propositional logic is a tool for reasoning about how various statements affect one another.
- To better understand how to prove a result, it often helps to translate what you're trying to prove into propositional logic first.
- That said, propositional logic isn't expressive enough to capture all statements. For that, we need something more powerful.

Next Time

- ***First-Order Logic***
 - Reasoning about groups of objects.
- ***First-Order Translations***
 - Expressing yourself in symbolic math!