

## Great Expectations

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Earlier in the course we came to the important result that  $E[\sum_i X_i] = \sum_i E[X_i]$ . First, as a warm up lets go back to our old friends and show how we could have derived expressions for their expectation.

### Expectation of Binomial

First let's start with some practice with the sum of expectations of indicator variables. Let  $Y \sim \text{Bin}(n, p)$ , in other words if  $Y$  is a Binomial random variable. We can express  $Y$  as the sum of  $n$  Bernoulli random indicator variables  $X_i \sim \text{Ber}(p)$ . Since  $X_i$  is a Bernoulli,  $E[X_i] = p$

$$Y = X_1 + X_2 + \dots + X_n = \sum_{i=1}^n X_i$$

Let's formally calculate the expectation of  $Y$ :

$$\begin{aligned} E[Y] &= E\left[\sum_i^n X_i\right] \\ &= \sum_i^n E[X_i] \\ &= E[X_0] + E[X_1] + \dots + E[X_n] \\ &= np \end{aligned}$$

### Expectation of Negative Binomial

Recall that a Negative Binomial is a random variable that semantically represents the number of trials until  $r$  successes. Let  $Y \sim \text{NegBin}(r, p)$ .

Let  $X_i = \#$  trials to get success after  $(i-1)$ st success. We can then think of each  $X_i$  as a Geometric RV:  $X_i \sim \text{Geo}(p)$ . Thus,  $E[X_i] = \frac{1}{p}$ . We can express  $Y$  as:

$$Y = X_1 + X_2 + \dots + X_r = \sum_{i=1}^r X_i$$

Let's formally calculate the expectation of  $Y$ :

$$\begin{aligned} E[Y] &= E\left[\sum_{i=1}^r X_i\right] \\ &= \sum_{i=1}^r E[X_i] \\ &= E[X_1] + E[X_2] + \dots + E[X_r] \\ &= \frac{r}{p} \end{aligned}$$

### Conditional Expectation

We have gotten to know a kind and gentle soul, conditional probability. And we now know another funky fool, expectation. Let's get those two crazy kids to play together.

Let  $X$  and  $Y$  be jointly random variables. Recall that the conditional probability mass function (if they are discrete), and the probability density function (if they are continuous) are respectively:

$$p_{X|Y}(x|y) = \frac{p_{X,Y}(x,y)}{p_Y(y)}$$

$$f_{X|Y}(x|y) = \frac{f_{X,Y}(x,y)}{f_Y(y)}$$

We define the conditional expectation of  $X$  given  $Y = y$  to be:

$$E[X|Y = y] = \sum_x x p_{X|Y}(x|y)$$

$$E[X|Y = y] = \int_{-\infty}^{\infty} x f_{X|Y}(x|y) dx$$

Where the first equation applies if  $X$  and  $Y$  are discrete and the second applies if they are continuous.

## Properties of Conditional Expectation

Here are some helpful, intuitive properties of conditional expectation:

$$E[g(X)|Y = y] = \sum_x g(x) p_{X|Y}(x|y) \quad \text{if X and Y are discrete}$$

$$E[g(X)|Y = y] = \int_{-\infty}^{\infty} g(x) f_{X|Y}(x|y) dx \quad \text{if X and Y are continuous}$$

$$E\left[\sum_{i=1}^n X_i | Y = y\right] = \sum_{i=1}^n E[X_i | Y = y]$$

$$E[E[X|Y]] = E[X]$$

### Example 1

You roll two 6-sided dice  $D_1$  and  $D_2$ . Let  $X = D_1 + D_2$  and let  $Y =$  the value of  $D_2$ . What is  $E[X|Y = 6]$

$$E[X|Y = 6] = \sum_x x P(X = x|Y = 6)$$

$$= \left(\frac{1}{6}\right) (7 + 8 + 9 + 10 + 11 + 12) = \frac{57}{6} = 9.5$$

Which makes intuitive sense since  $6 + E[\text{value of } D_1] = 6 + 3.5$

### Example 2

Consider the following code with random numbers:

```
int Recurse() {
    int x = randomInt(1, 3); // Equally likely values
    if (x == 1) return 3;
    else if (x == 2) return (5 + Recurse());
    else return (7 + Recurse());
}
```

Let  $Y =$  value returned by “Recurse”. What is  $E[Y]$ . In other words, what is the expected return value. Note that this is the exact same approach as calculating the expected run time.

$$E[Y] = E[Y|X = 1]P(X = 1) + E[Y|X = 2]P(X = 2) + E[Y|X = 3]P(X = 3)$$

First lets calculate each of the conditional expectations:

$$\begin{aligned} E[Y|X = 1] &= 3 \\ E[Y|X = 2] &= E[5 + Y] = 5 + E[Y] \\ E[Y|X = 3] &= E[7 + Y] = 7 + E[Y] \end{aligned}$$

Now we can plug those values into the equation. Note that the probability of X taking on 1, 2, or 3 is 1/3:

$$\begin{aligned} E[Y] &= E[Y|X = 1]P(X = 1) + E[Y|X = 2]P(X = 2) + E[Y|X = 3]P(X = 3) \\ &= 3(1/3) + (5 + E[Y])(1/3) + (7 + E[Y])(1/3) \\ &= 15 \end{aligned}$$

## Hiring Software Engineers

You are interviewing  $n$  software engineer candidates and will hire only 1 candidate. All orderings of candidates are equally likely. Right after each interview you must decide to hire or not hire. You can not go back on a decision. At any point in time you can know the relative ranking of the candidates you have already interviewed.

The strategy that we propose is that we interview the first  $k$  candidates and reject them all. Then you hire the next candidate that is better than all of the first  $k$  candidates. What is the probability that the best of all the  $n$  candidates is hired for a particular choice of  $k$ ? Let's denote that result  $P_k(\text{Best})$ . Let  $X$  be the position in the ordering of the best candidate:

$$\begin{aligned} P_k(\text{Best}) &= \sum_{i=1}^n P_k(\text{Best}|X = i)P(X = i) \\ &= \frac{1}{n} \sum_{i=1}^n P_k(\text{Best}|X = i) \end{aligned} \quad \text{since each position is equally likely}$$

What is  $P_k(\text{Best}|X = i)$ ? if  $i \leq k$  then the probability is 0 because the best candidate will be rejected without consideration. Sad times. Otherwise we will chose the best candidate, who is in position  $i$ , only if the best of the first  $i - 1$  candidates is among the first  $k$  interviewed. If the best among the first  $i - 1$  is not among the first  $k$ , that candidate will be chosen over the true best. Since all orderings are equally likely the probability that the best among the  $i - 1$  candidates is in the first  $k$  is:

$$\frac{k}{i-1} \quad \text{if } i > k$$

Now we can plug this back into our original equation:

$$\begin{aligned} P_k(\text{Best}) &= \frac{1}{n} \sum_{i=1}^n P_k(\text{Best}|X = i) \\ &= \frac{1}{n} \sum_{i=k+1}^n \frac{k}{i-1} \quad \text{since we know } P_k(\text{Best}|X = i) \\ &\approx \frac{1}{n} \int_{k+1}^n \frac{k}{i-1} di \quad \text{By Riemann Sum approximation} \\ &= \frac{k}{n} \ln(i-1) \Big|_{k+1}^n = \frac{k}{n} \ln \frac{n-1}{k} \approx \frac{k}{n} \ln \frac{n}{k} \end{aligned}$$

If we think of  $P_k(\text{Best}) = \frac{k}{n} \ln \frac{n}{k}$  as a function of  $k$  we can take find the value of  $k$  that optimizes it by taking its derivative and setting it equal to 0. The optimal value of  $k$  is  $n/e$ . Where  $e$  is Euler's number.