Fusion Trees
Part One
Outline for Today

• **Word-Level Parallelism**
  • Harnessing the intrinsic parallelism inside the processor.

• **Word-Parallel Operations**
  • Comparing, tiling, and ranking numbers; adding and packing bits.

• **The Sardine Tree**
  • Unconditionally beating a BST for very small integers.

• **Most-Significant Bits**
  • Finding the most significant bit in O(1) time/space.
Recap from Last Time
Our Machine Model

- We will assume we’re working on a machine where memory is segmented into $w$-bit words.
  - Although on any one fixed machine $w$ is a constant, in general, don’t assume this is the case. 32-bit was the norm until fairly recently, and before that 16-bit was standard.

- We’ll assume C integer operators work in constant time, and won’t assume other integer operations (say, finding most significant bits, counting 1 bits set) are available.

  $+\quad -\quad *\quad /\quad \%\quad <<\quad >>\quad &\quad |\quad ^\quad =\quad <=$
Integer Ordered Dictionaries

• The **y-Fast Trie** is an ordered dictionary for integers where each operation takes $O(\log w)$ time, amortized/expected.
  
  • Note that when $\log n = \omega(\log w)$, this is better than a binary search tree!

• Space usage is $\Theta(n)$ words, where $n$ is the number of elements in the trie.
New Stuff!
A Key Technique: *Word-Level Parallelism*
Word-Level Parallelism

- On a standard computer, arithmetic and logical operations on a machine word take time $O(1)$.

- We can perform certain classes of operations (addition, shifts, etc.) on $\Theta(w)$ bits in time $O(1)$.
  - Think of this as a weak form of parallel computation, where we can work over multiple bits in parallel with a limited set of operations.

- With some creativity, we can harness these primitives to build operations that run in time $O(1)$ but work on $\omega(1)$ objects.

- Let’s see a quick example...
Word-Level Parallelism

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Word-Level Parallelism

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Word-Level Parallelism

We’ve performed eight logical additions with a single add instruction!
The Landscape

- Preprocessing/runtime tradeoffs:
  
  “Yes, we have to do a lot of work, but it’s a one-time cost and everything is cheaper after that.”

- Randomization:
  
  “We might have to do a lot of work, but it’s unlikely that we’ll do so.”

- Amortization:
  
  “Yes, we have to do a lot of work every once and a while, but only after a period of doing very little.”

- Word-level parallelism:
  
  “We have to do a lot of work, but we don’t have to perform many operations to do it.”
These actually aren’t called sardine trees. I couldn’t find a name for them anywhere and thought that this title was appropriate. Let me know if there’s a more proper name to associate with them!
The Setup

• Let $w$ denote the machine word size.
• Imagine you want to store a collection of $s$-bit integers, where $s$ is small compared to $w$.
  • For example, storing 7-bit integers on a 64-bit machine would have $s = 7$ and $w = 64$.
• Can we build an ordered dictionary that takes advantage of the small key size?
A Refresher: B-Trees

- A **B-tree** is a multiway tree with a tunable parameter $b$ called the **order** of the tree.

- Each nodes stores $\Theta(b)$ keys. The height of the tree is $\Theta(\log_b n)$.

- Most operations (**lookup, insert, delete, successor, predecessor**, etc.) perform a top-down search of the tree, doing some amount of work per node.

- Runtime of each operation is $O(f(b) \log_b n)$, where $f(b)$ is the amount of work done per node.
B-Tree Traversals

- Most B-tree operations work by choosing some subtree to descend into, then descending there.

- **Claim:** The subtree we want is given by the number of keys in the current node less than or equal to the query key \( k \). This quantity is the **rank** of \( k \).

- For example, in the top node of the B-tree shown below:
  - \( \text{rank}(40) = 0 \)
  - \( \text{rank}(74) = 2 \)
  - \( \text{rank}(107) = 3 \)

- **Question:** How quickly can we determine the rank of a key in a B-tree node?
B-Tree Traversals

- We can determine \( \text{rank}(k) \) with a linear search in each B-tree node for a total lookup cost of \( O(b \cdot \log_b n) \).

- We can determine \( \text{rank}(k) \) with a binary search in each B-tree node for a total lookup cost of

\[
O(\log_b n \cdot \log b) = O(\log n).
\]

- **Claim:** If we can fit all the keys in a node into \( O(1) \) machine words, we can determine \( \text{rank}(k) \) in time \( O(1) \) for total lookup cost of \( O(\log_b n) \).
How is this possible?
Warmup: Comparing Two Values

- Imagine we have two $s$-bit integers $x$ and $y$ and want to determine whether $x \geq y$.
- How might we do this?

\[
\begin{array}{cccc}
0 & 1 & 1 & 0 \\
1 & 0 & 0 & 1 \\
\hline
0 & 0 & 1 & 1 \\
\hline
1 & 0 & 0 & 1 \\
\end{array}
\]

\[
\begin{array}{cccc}
? & 0 & 1 & 0 \\
1 & 1 & 0 & 0 \\
\hline
1 & 1 & 1 & 1 \\
\end{array}
\]
Warmup: Comparing Two Values

- Imagine we have two \( s \)-bit integers \( x \) and \( y \) and want to determine whether \( x \geq y \).
- How might we do this?

\[
\begin{array}{cccccccc}
1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\
- & 0 & 0 & 0 & 1 & 1 & 1 & 1 \\
\hline
1 & 1 & 0 & 0 & 1 & 1 & 0 & 1
\end{array}
\]
Warmup: Comparing Two Values

• Imagine we have two $s$-bit integers $x$ and $y$ and want to determine whether $x \geq y$.
• How might we do this?

This bit tells us whether the first number was as least as big as the second!
### Comparing Multiple Values

- This technique can be extended to work on multiple values in parallel.
- For example, here’s how we’d compare eight pairs of 7-bit numbers by doing a single 64-bit subtraction:

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Comparing Multiple Values

- This technique can be extended to work on multiple values in parallel.
- For example, here’s how we’d compare eight pairs of 7-bit numbers by doing a single 64-bit subtraction:

```
11011110 10101110 11110000 11001101 10101111 10001101 11101111 11000001
- 00011010 01001011 00010100 00100000 01010000 00100010 01000100 00001000
```

```
11010100 01101001 11100100 10101101 01011111 01101011 10110011 11011001
```
Comparing Multiple Values

• This technique can be extended to work on multiple values in parallel.

• For example, here’s how we’d compare eight pairs of 7-bit numbers by doing a single 64-bit subtraction:

```
11101110 10101110 11111000 11001101 10101111 10001101 1110111 11100001
- 00011010 0100101 00010100 00100000 01010000 00100010 0100010 00001000
```

```
11010100 01101001 11100100 10101101 01011111 01101011 10110011 11011001
```
Comparing Multiple Values

● This technique can be extended to work on multiple values in parallel.

● For example, here’s how we’d compare eight pairs of 7-bit numbers by doing a single 64-bit subtraction:

\[
\begin{array}{cccccccccccc}
1101110 & 10101110 & 11111000 & 11001101 & 10101111 & 10001101 & 1110111 & 1100001 \\
- 00011010 & 0100101 & 0010100 & 0100000 & 0101000 & 00100010 & 0100100 & 0001000 \\
\hline
11010100 & 01101001 & 11100100 & 10101101 & 01011111 & 01101011 & 10110011 & 11011001
\end{array}
\]

This technique is used in practice, including the glibc version of strlen. Thanks to former CS166 student Jane Lange for pointing this out!
Fundamental Primitive: *Parallel Compare*

**Input:** Two machine words. The first holds an array $x_1, \ldots, x_n$ with one bit of space between each number. The second holds an array $y_1, \ldots, y_n$ with one bit of space between each number.

**Output:** A machine word with the result of $x_i \geq y_i$ encoded as a bit in the blank spaces between the numbers in the input array.

**Procedure:**

1. Use a bitwise OR to place 1s between the $x_i$’s.
2. Use a bitwise AND to place 0s between the $y_i$’s.
3. Compute $X - Y$. The bit preceding $x_i - y_i$ is 1 if $x_i \geq y_i$ and 0 otherwise.
Back to B-Trees

• **Recall:** The whole reason we’re interested in making these comparisons is so that we can find how many keys in a B-tree node are less than or equal to a query key \( k \).

• **Idea:** Store the (s-bit) keys in the B-tree node in a single (w-bit) machine word, with zeros interspersed:

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<th>( y_1 )</th>
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<td>103</td>
<td>106</td>
<td>107</td>
<td>109</td>
<td>110</td>
<td>127</td>
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Back to B-Trees

• **Recall:** The whole reason we’re interested in making these comparisons is so that we can find how many keys in a B-tree node are less than or equal to a query key $k$.

• **Idea:** Store the (s-bit) keys in the B-tree node in a single (w-bit) machine word, with zeros interspersed:

\[
\begin{array}{cccccccc}
 y_1 & y_2 & y_3 & y_4 & y_5 & y_6 & y_7 & y_8 \\
00101001 & 01011101 & 01100111 & 01101010 & 01101111 & 01101111 & 01101110 & 01111111
\end{array}
\]
Rank in $O(1)$

- To perform a lookup for the key $k$, form a number by replicating $k$ multiple times with 1s interspersed.

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<td>$k$</td>
<td>1100111</td>
<td>1100111</td>
<td>01101001</td>
<td>01011101</td>
<td>01100111</td>
<td>01101010</td>
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Rank in O(1)

• To perform a lookup for the key $k$, form a number by replicating $k$ multiple times with 1s interspersed.
Rank in $O(1)$

- To perform a lookup for the key $k$, form a number by replicating $k$ multiple times with 1s interspersed.
- Subtract the B-tree key number from it to do a parallel comparison.

\[
\begin{array}{cccccccccc}
  k & k & k & k & k & k & k & k & k & k \\
  y_1 & y_2 & y_3 & y_4 & y_5 & y_6 & y_7 & y_8 \\
  00101001 & 01011101 & 01100111 & 01101010 & 01101011 & 01101101 & 01101110 & 01111111
\end{array}
\]
Rank in $O(1)$

- To perform a lookup for the key $k$, form a number by replicating $k$ multiple times with 1s interspersed.
- Subtract the B-tree key number from it to do a parallel comparison.

```
1100111 1100111 1100111 1100111 1100111 1100111 1100111 1100111
- 00101001 01011101 01100111 01101010 01101011 01101101 01101110 01111111
```

```
10111110 10001010 10000000 01111101 01111100 01111010 01111001 01101000
```
Rank in O(1)

- To perform a lookup for the key $k$, form a number by replicating $k$ multiple times with 1s interspersed.
- Subtract the B-tree key number from it to do a parallel comparison.
- Count up how many of the sentinel bits in the resulting number are equal to 1. This is the number of keys in the node less than or equal to $k$.

```
1100111 1100111 1100111 1100111 1100111 1100111 1100111 1100111
- 00101001 01011101 01100111 01101010 01101011 01101101 01101110 01111111
```

Rank: 3
Back in Base Ten

- Suppose you have a one-digit number $m$.
- You want to form this base-10 number:
  \[mmm\]
- Is there a nice series of arithmetical operations that will produce this?
- **Answer:** Compute $m \times 111$.
- Why does this work?
  \[
m \times 111 = m \times 100 + m \times 10 + m \times 1
  = m00 + 0m0 + 00m
  = mmm.
  \]
Back in Base Ten

- Suppose you have a one-digit number \( m \).
- You want to form this base-10 number: \( mmm \).
- Is there a nice series of arithmetical operations that will produce this?
- \textbf{Answer:} Compute \( m \times 111 \).
- Why does this work?

\[
m \times 111 = m \ll 2 + m \ll 1 + m \ll 0 = m00 + 0m0 + 00m = mmm.
\]
Back in Base Ten

• Suppose you have a two-digit number $mn$.
• You want to form this base-10 number:

$$mnmnmnmn$$

• Is there a nice series of arithmetical operations that will produce this?

**Answer:** Compute $mn \times 10,101$.

• Why does this work?

\[
mn \times 10,101 = mn \times 10,000 + mn \times 100 + mn \times 1
\]
\[
= mn0000 + 00mn00 + 0000mn
\]
\[
= mnmnmnmn.
\]
Back in Base Ten

• Suppose you have a two-digit number $mn$.
• You want to form this base-10 number:

$$mnmnmn$$

• Is there a nice series of arithmetical operations that will produce this?
• **Answer:** Compute $mn \times 10,101$.
• Why does this work?

$$mn \times 10,101 = mn \ll 4 + mn \ll 2 + mn \ll 0$$

$$= mn0000 + 00mn00 + 0000mn$$

$$= mnmnmn.$$
Back in Base Ten

**Answer:** Compute $mn \times 10,101$.

Why does this work?

$$mn \times 10,101 = mn \ll 4 + mn \ll 2 + mn \ll 0$$

$$= mn0000 + 00mn00 + 0000mn$$

$$= mnmnmn.$$

Back in Base Ten

• Suppose you have a three-digit number $mnp$.
• You want to form this base-10 number:
  
  $mnp000mnp0mnp$

• Is there a nice series of arithmetical operations that will produce this?

• **Answer:** Compute $mnp \times /*$ something */.
  
  $mnp000mnp0mnp$

  $= mnp \ll 10 + mnp \ll 4 + mnp \ll 0$

  $= mnp \times 10^{10} + mnp \times 10^{4} + mnp \times 10^{0}$

  $= mnp \times 10,000,010,001$
Back in Base Ten

• Suppose you have a three-digit number $mnp$.
• You want to form this base-10 number:
  $$mnp000mnp0mnp$$
• Is there a nice series of arithmetical operations that will produce this?
• **Answer:** Compute $mnp \times 10,000,010,001$.
  $$mnp000mnp0mnp$$
  $$= mnp \ll 10 + mnp \ll 4 + mnp \ll 0$$
  $$= mnp \times 10^{10} + mnp \times 10^4 + mnp \times 10^0$$
  $$= mnp \times 10,000,010,001$$
Computing Rank in O(1)

```c
const uint64_t kMultiplier = 0b1000000010000000...100000001;

uint64_t tiledK = k * kMultiplier;
```
Computing Rank in O(1)

```c
const uint64_t kMultiplier = 0b1000000010000000...100000001;

uint64_t tiledK = k * kMultiplier;
```

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<td>1100111</td>
<td>1100111</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>y1</th>
<th>y2</th>
<th>y3</th>
<th>y4</th>
<th>y5</th>
<th>y6</th>
<th>y7</th>
<th>y8</th>
</tr>
</thead>
<tbody>
<tr>
<td>00101001</td>
<td>01011101</td>
<td>01100111</td>
<td>01101010</td>
<td>01101011</td>
<td>01101101</td>
<td>01101110</td>
<td>01111111</td>
</tr>
</tbody>
</table>
Fundamental Primitive: *Parallel Tile*

**Input:** A number $k$ much smaller than a machine word.

**Output:** A machine word holding multiple tiled copies of $k$, spread out with gaps between each copy.

**Procedure:**

1. Form a number $M$ with a 1 bit at the end of each location to tile $k$.
2. Compute $M \times k$. 
Computing Rank in O(1)

```cpp
const uint64_t kMultiplier = 0b1000000010000000...100000001;

uint64_t tiledK = k * kMultiplier;
```
Computing Rank in O(1)

```c
const uint64_t kMultiplier = 0b1000000010000000...100000001;
const uint64_t kOnesMask   = 0b1000000010000000...010000000;

uint64_t tiledK = k * kMultiplier;
```

<table>
<thead>
<tr>
<th>$k$</th>
<th>$k$</th>
<th>$k$</th>
<th>$k$</th>
<th>$k$</th>
<th>$k$</th>
<th>$k$</th>
<th>$k$</th>
</tr>
</thead>
<tbody>
<tr>
<td>1100111</td>
<td>1100111</td>
<td>1100111</td>
<td>1100111</td>
<td>1100111</td>
<td>1100111</td>
<td>1100111</td>
<td>1100111</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>$y_1$</th>
<th>$y_2$</th>
<th>$y_3$</th>
<th>$y_4$</th>
<th>$y_5$</th>
<th>$y_6$</th>
<th>$y_7$</th>
<th>$y_8$</th>
</tr>
</thead>
<tbody>
<tr>
<td>00101001</td>
<td>01011101</td>
<td>01100111</td>
<td>01101010</td>
<td>01101011</td>
<td>01101101</td>
<td>01101110</td>
<td>01111111</td>
</tr>
</tbody>
</table>
Computing Rank in $O(1)$

```cpp
const uint64_t kMultiplier = 0b1000000010000000...100000001;
const uint64_t kOnesMask   = 0b1000000001000000...01000000;

uint64_t tiledK = (k * kMultiplier) | kOnesMask;
```
Computing Rank in O(1)

```c
const uint64_t kMultiplier = 0b1000000010000000...100000001;
const uint64_t kOnesMask   = 0b1000000010000000...010000000;

uint64_t tiledK     = (k * kMultiplier) | kOnesMask;
```

```plaintext
11100111  11100111  11100111  11100111  11100111  11100111  11100111  11100111  
00101001  01011101  01100111  01101010  01101011  01101101  01101110  01111111
```
Computing Rank in $O(1)$

```c
const uint64_t kMultiplier = 0b1000000010000000...100000001;
const uint64_t kOnesMask = 0b1000000010000000...0100000001;
uint64_t tiledK = (k * kMultiplier) | kOnesMask;
uint64_t comparison = tiledK - packedKeys;
```
Computing Rank in $O(1)$

```cpp
const uint64_t kMultiplier = 0b1000000010000000...100000001;
const uint64_t kOnesMask   = 0b1000000010000000...010000000;

uint64_t tiledK     = (k * kMultiplier) | kOnesMask;
uint64_t comparison = (tiledK - packedKeys) & kOnesMask;
```

How do we count how many of these bits are set?
Summing Up Flags

• After performing our subtraction, we’re left with a number like this one, where the highlighted bits are “interesting” to us.

• **Goal:** Add up these “interesting” values using at most O(1) total operations on words.
An Initial Idea

• To sum up the flags, we could extract each bit individually and add the result.

• **The catch:** This takes time $\Theta(r)$, where $r$ is the number of times we tiled our value.

• Can we do better?

```
 a00000000 b00000000 c00000000 d00000000
```
A Shifty Solution

• Given this number:

\[ \begin{array}{cccc}
  a & b & c & d \\
 0 & 0 & 0 & 0 \\
\end{array} \]

we want to compute \( a + b + c + d \).

• We can’t efficiently isolate \( a, b, c, \) and \( d \).

• **Claim:** We don’t have to!

\[ \begin{array}{cccc}
  d \\
  c \\
  b \\
  a \\
\end{array} \]

\[ \begin{array}{cccc}
  d \\
  c \\
  b \\
  a \\
\end{array} \]

\[ \hline \]

\[ \begin{array}{cccc}
  d \\
  c \\
  b \\
  a \\
\end{array} \]

\[ \hline \]

\[ \begin{array}{cccc}
  d \\
  c \\
  b \\
  a \\
\end{array} \]
A Shifty Solution

• Given this number:
  \[ \begin{array}{cccc}
    a & 0 \cdots & 0 & a \\
    b & 0 \cdots & 0 & b \\
    c & 0 \cdots & 0 & c \\
    d & 0 \cdots & 0 & d \\
  \end{array} \]

we want to compute \( a + b + c + d \).

• We can’t efficiently isolate \( a, b, c, \) and \( d \).

• **Claim:** We don’t have to!

\[ \begin{array}{cccc}
    d & + & b & + \\
    c & + & a & + \\
    \hline
    & & & \end{array} \]
A Shifty Solution

• Given this number:

\[
\begin{array}{cccc}
a & b & c & d \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
\end{array}
\]

we want to compute \( a + b + c + d \).

• We can’t efficiently isolate \( a, b, c, \) and \( d \).

• **Claim:** We don’t have to!

\[
\begin{array}{cccc}
d \\
c \\
b \\
0 \\
0 \\
0 \\
\end{array}
\]

\[
\begin{array}{cccc}
a & b & c & d \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
+ \\
\end{array}
\]

\[
\begin{array}{cccc}
a & b & c & d \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
\end{array}
\]
A Shifty Solution

• Given this number:

\[
\begin{align*}
a & \quad b & \quad c & \quad d \\
0000000 & \quad 0000000 & \quad 0000000 & \quad 0000000
\end{align*}
\]

we want to compute \( a + b + c + d \).

• We can’t efficiently isolate \( a, b, c, \) and \( d \).

• **Claim:** We don’t have to!

\[
\begin{align*}
da & \quad c \\
0000000 & \quad 0000000 & \quad 0000000 & \quad 0000000
\end{align*}
\]

\[
\begin{align*}
+ & \quad a & \quad b & \quad c & \quad d \\
0000000 & \quad 0000000 & \quad 0000000 & \quad 0000000
\end{align*}
\]

\[
\begin{align*}
\hline
& \quad a & \quad b & \quad c & \quad d \\
0000000 & \quad 0000000 & \quad 0000000 & \quad 0000000
\end{align*}
\]
A Shifty Solution

• Given this number:

\[ \begin{array}{cccc}
\text{a} & \text{b} & \text{c} & \text{d} \\
0000000 & 0000000 & 0000000 & 0000000 \\
\end{array} \]

we want to compute \( \text{a} + \text{b} + \text{c} + \text{d} \).

• We can’t efficiently isolate \( \text{a} \), \( \text{b} \), \( \text{c} \), and \( \text{d} \).

• Claim: We don’t have to!

\[ \begin{array}{cccc}
\text{d} \\
0000000 & 0000000 & 0000000 & 0000000 \\
\text{a} & \text{b} & \text{c} & \text{d} \\
0000000 & 0000000 & 0000000 & 0000000 \\
\text{a} & \text{b} & \text{c} & \text{d} \\
0000000 & 0000000 & 0000000 & 0000000 \\
\end{array} \]

\[ \begin{array}{cccc}
\text{+} \\
\text{a} & \text{b} & \text{c} & \text{d} \\
0000000 & 0000000 & 0000000 & 0000000 \\
\end{array} \]

\[ \begin{array}{cccc}
\text{a} & \text{b} & \text{c} & \text{d} \\
0000000 & 0000000 & 0000000 & 0000000 \\
\end{array} \]
A Shifty Solution

• Given this number:

\[ \begin{array}{cccc}
  a & b & c & d \\
 0 & 0 & 0 & 0 \\
\end{array} \]

we want to compute \( a + b + c + d \).

• We can’t efficiently isolate \( a, b, c, \) and \( d \).

• **Claim:** We don’t have to!

\[ \begin{array}{cccc}
  a & b & c & d \\
 0 & 0 & 0 & 0 \\
\end{array} +
\]

\[ \begin{array}{cccc}
  a & b & c & d \\
 0 & 0 & 0 & 0 \\
\end{array} \]

\[ \begin{array}{cccc}
  a & b & c & d \\
 0 & 0 & 0 & 0 \\
\end{array} +
\]

\[ \begin{array}{cccc}
  a & b & c & d \\
 0 & 0 & 0 & 0 \\
\end{array} \]

\[ \begin{array}{cccc}
  a & b & c & d \\
 0 & 0 & 0 & 0 \\
\end{array} \]
A Shifty Solution

• Given this number:

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

we want to compute \(a + b + c + d\).

• We can’t efficiently isolate \(a\), \(b\), \(c\), and \(d\).

• Claim: We don’t have to!

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]

\[
\begin{array}{cccc}
 a & b & c & d \\
0 & 0 & 0 & 0 \\
\end{array}
\]
A Shifty Solution

• Given this number:

\[
\begin{array}{cccc}
  a & 0 & 0 & 0 \\
  b & 0 & 0 & 0 \\
  c & 0 & 0 & 0 \\
  d & 0 & 0 & 0 \\
\end{array}
\]

we want to compute \( a + b + c + d \).

• We can’t efficiently isolate \( a \), \( b \), \( c \), and \( d \).

• **Claim:** We don’t have to!

\[
\begin{array}{cccc}
  a & 0 & 0 & 0 \\
  b & 0 & 0 & 0 \\
  c & 0 & 0 & 0 \\
  d & 0 & 0 & 0 \\
\end{array}
+ \quad \begin{array}{cccc}
  a & 0 & 0 & 0 \\
  b & 0 & 0 & 0 \\
  c & 0 & 0 & 0 \\
  d & 0 & 0 & 0 \\
\end{array}
= \quad \begin{array}{cccc}
  a & 0 & 0 & 0 \\
  b & 0 & 0 & 0 \\
  c & 0 & 0 & 0 \\
  d & 0 & 0 & 0 \\
\end{array}
\]

This is a series of shifts and adds. It’s equivalent to multiplying our original number by some well-chosen spreader!
Fundamental Primitive: *Parallel Add*

**Input:** A machine word with “interesting” bits spaced evenly across the word.

**Output:** The sum of those “interesting” bits.

**Procedure:**

1. Perform a *parallel tile* with an appropriate multiplier to place all leading bits on top of one another.

2. Use a bitmask and bitshift to isolate those bits.
Computing Rank in O(1)

```c
const uint64_t kMultiplier = 0b1000000010000000...100000001;
const uint64_t kOnesMask   = 0b1000000010000000...010000000;

uint64_t tiledK     = (k * kMultiplier) | kOnesMask;
uint64_t comparison = (tiledK – packedKeys) & kOnesMask;
```

a0000000 b0000000 c0000000 d0000000
Computing Rank in O(1)

```c
const uint64_t kMultiplier = 0b1000000010000000...100000001;
const uint64_t kOnesMask   = 0b1000000010000000...010000000;

uint64_t tiledK     = (k * kMultiplier) | kOnesMask;
uint64_t comparison = (tiledK – packedKeys) & kOnesMask;

const uint64_t kStacker = 0b1000001000001...1000001;
const uint8_t  kShift   = 31;
const uint64_t kMask    = 0b111;

uint64_t rank = comparison * kStacker;
```

a0000000 b0000000 c0000000 d0000000
Computing Rank in O(1)

```c
const uint64_t kMultiplier = 0b1000000010000000...100000001;
const uint64_t kOnesMask = 0b1000000010000000...010000000;

uint64_t tiledK = (k * kMultiplier) | kOnesMask;
uint64_t comparison = (tiledK - packedKeys) & kOnesMask;

const uint64_t kStacker = 0b1000001000001...1000001;
const uint8_t kShift = 31;
const uint64_t kMask = 0b111;

uint64_t rank = comparison * kStacker;
```

Computing Rank in $O(1)$

```c
const uint64_t kMultiplier = 0b1000000010000000...100000001;
const uint64_t kOnesMask = 0b1000000010000000...010000000;
uint64_t tiledK = (k * kMultiplier) | kOnesMask;
uint64_t comparison = (tiledK - packedKeys) & kOnesMask;
const uint64_t kStacker = 0b1000001000001...1000001;

const uint8_t kShift = 31;
const uint64_t kMask = 0b111;
uint64_t rank = comparison * kStacker;
```

```
+  a00000000 b00000000 c00000000 d00000000 00000000 00000000 00000000 00000000 00000000
  a00000000 b00000000 c00000000 d00000000 00000000 00000000 00000000 00000000 00000000
  a00000000 b00000000 c00000000 d00000000 00000000 00000000 00000000 00000000 00000000
  a00000000 b00000000 c00000000 d00000000 00000000 00000000 00000000 00000000 00000000
  a00000000 b00000000 c00000000 d00000000 00000000 00000000 00000000 00000000 00000000

???????? ????sum???????? ?? ????????? ???? ????????? ???? ????
Computing Rank in O(1)

```c
const uint64_t kMultiplier = 0b1000000010000000...100000001;
const uint64_t kOnesMask   = 0b1000000010000000...010000000;
uint64_t tiledK     = (k * kMultiplier) | kOnesMask;
uint64_t comparison = (tiledK – packedKeys) & kOnesMask;
const uint64_t kStacker = 0b1000001000001...1000001;
const uint8_t  kShift   = 31;
uint64_t rank =   comparison * kStacker;
```

```
a0000000 b0000000 c0000000 d0000000 00000000 00000000 00000000 00000000
  a0000000 b0000000 c0000000 d0000000 00000000 00000000 00000000 00000000
    a0000000 b0000000 c0000000 d0000000 00000000 00000000 00000000 00000000
+  a0000000 b0000000 c0000000 d0000000 00000000 00000000 00000000 00000000

???????? ????????? ?????? sum???????? ????????? ????????? ????????? ?????
Computing Rank in O(1)

```cpp
const uint64_t kMultiplier = 0b1000000010000000...100000001;
const uint64_t kOnesMask   = 0b1000000010000000...010000000;
uint64_t tiledK     = (k * kMultiplier) | kOnesMask;
uint64_t comparison = (tiledK – packedKeys) & kOnesMask;

const uint64_t kStacker = 0b1000001000001...1000001;
const uint8_t  kShift   = 31;

uint64_t rank =  (comparison * kStacker) >> kShift;
```

```
a0000000 b0000000 c0000000 d0000000 00000000 00000000 00000000 00000000

a0000000 b0000000 c0000000 d0000000 00000000 00000000 00000000 00000000

a0000000 b0000000 c0000000 d0000000 00000000 00000000 00000000 00000000

+  

a0000000 b0000000 c0000000 d0000000 00000000 00000000 00000000 00000000

???????? ????????? ????????? ????????? ????????? ????????? ????????? sum
```
Computing Rank in O(1)

\[
\text{const} \ \ \text{uint64}_t \ \ k\text{Multiplier} = 0b1000000010000000\ldots100000001; \\
\text{const} \ \ \text{uint64}_t \ \ k\text{OnesMask} = 0b1000000010000000\ldots01000000;
\]

\[
\text{uint64}_t \ \ \text{tiledK} = (k \times k\text{Multiplier}) | k\text{OnesMask}; \\
\text{uint64}_t \ \ \text{comparison} = (\text{tiledK} - \text{packedKeys}) \& k\text{OnesMask}; \\
\text{const} \ \ \text{uint64}_t \ \ k\text{Stacker} = 0b1000001000001\ldots1000001; \\
\text{const} \ \ \text{uint8}_t \ \ k\text{Shift} = 31; \\
\text{const} \ \ \text{uint64}_t \ \ k\text{Mask} = 0b111;
\]

\[
\text{uint64}_t \ \ \text{rank} = (\text{comparison} \times k\text{Stacker}) \gg k\text{Shift};
\]
Computing Rank in O(1)

```c
const uint64_t kMultiplier = 0b1000000010000000...100000001;
const uint64_t kOnesMask   = 0b1000000010000000...010000000;

uint64_t tiledK     = (k * kMultiplier) | kOnesMask;
uint64_t comparison = (tiledK - packedKeys) & kOnesMask;

const uint64_t kStacker = 0b1000001000001...1000001;
const uint8_t  kShift   = 31;
const uint64_t kMask    = 0b111;

uint64_t rank = ((comparison * kStacker) >> kShift) & kMask;
```

```
+  a0000000 b0000000 c0000000 d0000000
  a0000000 b0000000 c0000000 d0000000
  a0000000 b0000000 c0000000 d0000000
  a0000000 b0000000 c0000000 d0000000
  a0000000 b0000000 c0000000 d0000000

sum
```
Fundamental Primitive: *Parallel Rank*

**Input:** An array of integers packed into a machine word with one bit of space between integers, and a key \( k \).

**Output:** How many elements of the array are less than or equal to \( k \).

**Procedure:**

1. Perform a *parallel tile* to create \( n \) copies of the key \( k \), prefixed by 1’s.

2. Perform a *parallel compare* of the key \( k \) against values \( x_1, \ldots, x_n \).

3. Perform a *parallel add* to sum those values into some total \( t \).

4. Return \( t \).
The Sardine Tree

- Let $w$ be the word size and $s$ be some (much) smaller number of bits.
- A **sardine tree** is a B-tree of order $\Theta(w/s)$ where the keys in a node are packed into a single machine word.
  - Get it? The keys are “packed” tightly into a machine word! I’m funny.
- Each node is annotated with several values (the masks and multipliers from the preceding slide), which are updated in time $O(1)$ whenever a key is added or removed.
- Supports all ordered dictionary operations in time $O(\log_b n) = O(\log_{w/s} n)$. 
Here’s the performance breakdown for the sardine tree.

Notice that the runtime performance is strictly better than that of a BST!

Notice that the space usage is sublinear, since each node stores multiple keys!

The Sardine Tree

- **lookup**: $O(\log_{w/s} n)$
- **insert**: $O(\log_{w/s} n)$
- **delete**: $O(\log_{w/s} n)$
- **max**: $O(\log_{w/s} n)$
- **succ**: $O(\log_{w/s} n)$
- Space: $\Theta(n \cdot s/w)$ words
For Comparison

- It’s helpful to compare the sardine tree against the y-fast trie.
- To make the comparison fair, I’m treating the y-fast trie as though the machine word size is \( s \) rather than \( w \).

The y-Fast Trie

- \textbf{lookup}: \( O(\log s) \)
- \textbf{insert}: \( O(\log s) \)\(^*\)
- \textbf{delete}: \( O(\log s) \)\(^*\)
- \textbf{max}: \( O(\log s) \)
- \textbf{succ}: \( O(\log s) \)
- Space: \( \Theta(n) \) words

\(^*\) Expected, amortized
What’s Next

- **Question**: Can we get performance along these lines even if the keys fill full machine words?
- The strategy used in the sardine tree on its own won’t get us there – but many of those same techniques will!
- We’ll see how to do this next time. In the meantime, let’s see some other cool tricks we can do with word-level parallelism.

**Mystery Structure?**

- **lookup**: $O(\log_w n)$
- **insert**: $O(\log_w n)$
- **delete**: $O(\log_w n)$
- **max**: $O(\log_w n)$
- **succ**: $O(\log_w n)$
- Space: $\Theta(n)$
Word-Level Parallelism Tricks #2: Most-Significant Bits
Most-Significant Bits

• The *most-significant bit* function, denoted $\text{msb}(n)$, outputs the index of the highest 1 bit set in the binary representation of number $n$.

• Some examples:
  
  $\text{msb}(0110) = 2$  
  $\text{msb}(010100) = 4$  
  $\text{msb}(1111) = 3$

• Note that $\text{msb}(0)$ is undefined.

• Mathematically, $\text{msb}(n)$ is the largest value of $k$ such that $2^k \leq n$. (*Do you see why?*)
Most-Significant Bits

• Although we didn’t have this name earlier in the quarter, you’ve seen a place where we needed to efficiently compute $\text{msb}(n)$.

• Do you remember where?

• **Answer:** In the sparse table RMQ structure, where computing $\text{RMQ}(i, j)$ requires computing the largest number $k$ where $2^k \leq j - i + 1$.

• That’s exactly the value of $\text{msb}(j - i + 1)$!
Most-Significant Bits

• On many architectures, there’s a single assembly instruction that computes $\text{msb}(n)$.
  • on x86, it’s BSR (bit scan reverse).
• On others, nothing like this exists.
  • Older versions of MIPS, for example.
• **Question:** How would we compute $\text{msb}(n)$ assuming we only have access to the regular C operators?
  
  $$+ - \times \div \% \ll \gg \& | \^ == <=$$
Computing $\text{msb}$

- In Problem Set 1, you (probably) computed $\text{msb}(n)$ by building a lookup table mapping each value of $n$ to $\text{msb}(n)$.

- **The Good**: This takes time $O(1)$ to evaluate.

- **The Bad**: The preprocessing time, and space usage, is $\Theta(U)$, where $U$ is the maximum value we’ll be querying for.

- **The Ugly**: In the worst case $U = 2^w$.

- Can we do better?
Most-Significant Digits

- Can you compute most-significant digits
  - ... in time $O(w)$ using $O(1)$ space?
  - ... in time $O(\log w)$ using $O(1)$ space?
  - ... in time $O(1)$ using $O(1)$ space?
- Remember that the word size $w$ is not a constant and that we can only use C-style operations.
Most-Significant Digits

Can you compute most-significant digits

- ... in time $O(w)$ using $O(1)$ space?
  ... in time $O(\log w)$ using $O(1)$ space?
  ... in time $O(1)$ using $O(1)$ space?

Remember that the word size $w$ is not a constant and that we can only use C-style operations.
Most-Significant Bits

- There’s a simple $O(w)$-time algorithm for computing $\text{msb}(n)$ that just checks all the bits until a 1 is found:

```c
for (uint8_t bit = 64; bit > 0; bit--) {
    if (n & (uint64_t(1) << (bit - 1))) {
        return bit;
    }
}
flailAndPanic();
```

- Can we do better?
Computing msb

• We can improve this runtime to O(log \(w\)) by using a binary search:
  • Check if any bits in the upper half of the bits of \(n\) are set.
  • If so, recursively explore the upper half of \(n\).
  • If not, recursively explore the lower half of \(n\).

• We can test whether any bit in a range is set by ANDing with a mask of 1s and seeing if the result is nonzero:

  \[
  \begin{array}{cccccccccccc}
  11011100 & 10111011 & 11000100 & 11010101 & 11100110 & 11110111 & 11000010 & 00110010 \\
  \wedge & 11111111 & 11111111 & 11111111 & 11111111 & 00000000 & 00000000 & 00000000 & 00000000 \\
  \end{array}
  \]

• Can we do better?
Claim: For any machine word size $w$, there is an algorithm that uses $O(1)$ machine operations and $O(1)$ space – independently of $w$ – and computes $\text{msb}(n)$. This is not obvious!
How is this possible?
Not Starting from Scratch

• We’re not going into this problem blind. We’ve seen a bunch of useful techniques so far:
  • \textbf{Parallel compare:} We can compare a bunch of small numbers in parallel in $O(1)$ machine word operations.
  • \textbf{Parallel tile:} We can take a small number and “tile” it multiple times in $O(1)$ machine word operations.
  • \textbf{Parallel add:} If we have a bunch of “flag” bits spread out evenly, we can add them all up in $O(1)$ machine word operations.
  • \textbf{Parallel rank:} We can find the rank of a small number in an array of small numbers in $O(1)$ machine word operations.
• This is an impressive array of techniques. Let’s see if we can reuse or adapt them.
Not Starting from Scratch

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  • **Parallel compare**: We can compare a bunch of small numbers in parallel in $O(1)$ machine word operations.
  
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  • **Parallel rank**: We can find the rank of a small number in an array of small numbers in $O(1)$ machine word operations.
  
• This is an impressive array of techniques. Let’s see if we can reuse or adapt them.
MSBs as Ranks

- **Recall:** \( \text{msb}(n) \) is the largest value of \( k \) for which \( 2^k \leq n \).

- **Idea:** Imagine we have an array of all the powers of two that we can represent in a machine word. Then \( \text{msb}(n) \) is the rank of \( n \) in that array!

\[
\begin{array}{ccccccccccc}
2^0 & 2^1 & 2^2 & 2^3 & 2^4 & 2^5 & 2^6 & \cdots & 2^{50} & 2^{51} & 2^{52} & 2^{53} & 2^{54} & 2^{55} & 2^{56} & 2^{57} & 2^{58} & \cdots & 2^{61} & 2^{62} & 2^{63}
\end{array}
\]
The Problem

• We can compute the rank of a value in an array assuming that all the array entries fit into a single machine word.

• This isn’t the case here:
  • \( w \) total powers of two to write out.
  • Total bits needed: \( \Theta(w^2) \), way too big to fit into a word.

• **Question:** Can we still harness the benefits of this parallel rank operation?
A Nice Decomposition

- Imagine we want to compute the most-significant bit of a $w$-bit integer.
  - In what follows, we’ll pick $w = 64$, but this works for any $w$.
- We ultimately want to be finding the MSB of numbers with way fewer than $w$ bits.
- **Idea:** Split $w$ into some number of blocks of size $b$. Then,
  - find the index of the highest block with at least one 1 bit set,
  - find the index of the highest bit within that block.
Imagine we want to compute the most-significant bit of a \( w \)-bit integer.

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- find the index of the highest block with at least one 1 bit set, then
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Compute \( \text{msb} \) for a \( \frac{w}{b} \)-bit number.
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A Nice Decomposition

• We will compute the MSB for $w$-bit integers by solving MSB for $b$ and $w/b$-bit integers.

• What choice of $b$ minimizes $\max\{b, \frac{w}{b}\}$?

• **Answer:** Pick $b = \sqrt{w}$.

• So now we need to see how to
  
  • solve $\text{msb}(n)$ for integers with $\sqrt{w}$ bits, and
  
  • replace each block with a bit indicating whether that block contains a 1.
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- So now we need to see how to
  - solve $\text{msb}(n)$ for integers with $w^{1/2}$ bits, and
  - replace each block with a bit indicating whether that block contains a 1.
MSB for $w^{1/2}$ Bits

- **Recall**: We can compute $\text{msb}(n)$ by counting how many powers of two are less than or equal to $n$.
- If our numbers have size $w^{1/2}$, there are $w^{1/2}$ powers of two to compare against.
- Each of those powers of two has $w^{1/2}$ bits, so all of those powers of two can be packed into a single machine word!
- **Idea**: Use our O(1)-time rank algorithm!
MSB for $w^{1/2}$ Bits

- If our numbers have size $w^{1/2}$, there are $w^{1/2}$ powers of two to compare against, each of which has $w^{1/2}$ bits.
- Our parallel comparison prepends an extra bit to each number to compare.
MSB for $w^{1/2}$ Bits

- If our numbers have size $w^{1/2}$, there are $w^{1/2}$ powers of two to compare against, each of which has $w^{1/2}$ bits.
- Our parallel comparison prepends an extra bit to each number to compare.
- That’s barely – just barely – too many bits to fit into a machine word.
MSB for $w^{\frac{1}{2}}$ Bits

- **Claim:** This is an engineering problem at this point.
- **Option 1:** Split the powers of two into two different machine words and do two rank calculations.
- **Option 2:** Special-case the most-significant bit to reduce the number of bits to check.
- Either way, we find that the work done here is $O(1)$ machine operations, with no dependency on the word size $w$!
A Nice Decomposition

- We need to see how to
  - solve $\text{msb}(n)$ for integers with $w^{\frac{1}{2}}$ bits, and
  - replace each block with a bit indicating whether that block contains a 1.
A Nice Decomposition

- We need to see how to
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A Nice Decomposition

- We need to see how to
  - solve $\text{msb}(n)$ for integers with $\lceil \sqrt{n} \rceil$ bits, and
  - replace each block with a bit indicating whether that block contains a 1.
### Identifying Active Blocks

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**Observation:** A block contains a 1 bit if its first bit is 1 or its lower 7 bits contain a 1.
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Identifying Active Blocks

A number’s lower 7 bits contain a 1 if and only if the numeric value of those bits is at least 1.

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**High bit set?**

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High bit set?
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---

**Low bits set?**

| 00000000 | 10000000 | 10000000 | 10000000 | 00000001 | 10000000 | 10000000 | 10000000 |

---

**High bit set?**

| 00000000 | 00000000 | 10000000 | 10000000 | 00000000 | 10000000 | 10000000 | 10000000 |
Identifying Active Blocks

**Observation:** A block contains a 1 bit if its first bit is 1 or its lower 7 bits contain a 1.
Identifying Active Blocks

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Identifying Active Blocks

- We now have a word holding flags telling us which blocks have a 1 bit set.
- We need to find the highest set flag.
- There are only $w^{1/2}$ flags. If we could compact them into $w^{1/2}$ adjacent bits, we could use our earlier algorithm to find the highest one set!
Identifying Active Blocks

- **Idea:** Adapt the shifting technique we used to compute ranks.
- Instead of shifting the bits on top of one another, shift the bits next to one another:

```
  a\hline b\hline c\hline d
  \hline
  a \hline b \hline c \hline d
```

Identifying Active Blocks

- **Idea:** Adapt the shifting technique we used to compute ranks.

- Instead of shifting the bits on top of one another, shift the bits next to one another:

  \[ \begin{align*}
  &a00000000b00000000c00000000d00000000 \\
  &a00000000b00000000c00000000d00000000 \\
  &a00000000b00000000c00000000d00000000 \\
  \textbf{+} &a00000000b00000000c00000000d00000000 \\
  \textbf{?} &a00000000b00000000c00000000d00000000
  \end{align*} \]
Fundamental Primitive: **Parallel Pack**

**Input:** A machine word containing several “interesting” bits that are evenly spaced apart.

**Output:** A machine word with those “interesting” bits placed adjacent to one another at the low end of the word.

**Procedure:**

1. Perform a *parallel tile* with an appropriate multiplier to place all leading bits adjacent to one another.

2. Use a bitmask and bitshift to isolate those bits.
Putting It All Together

- Use a bitmask to identify all blocks whose high bit is set.
- Use a *parallel tile* and a *parallel compare* to identify all blocks with a 1 bit aside from the first.
- Use a *parallel pack* to pack those bits together.
- Use a *parallel rank* to determine the highest of those bits set, which gives the block index.
- Use a *parallel rank* to determine the highest bit set within that block.
The Finished Product

• I’ve posted a link to a working implementation of this algorithm for 64-bit integers on the course website.

• Feel free to check it out – it’s really magical seeing all the techniques come together!
What We Covered

- We can use bit-parallel tricks to
  - compare multiple values in parallel,
  - tile a number across a word,
  - sum up evenly-spaced bits in a word,
  - compute ranks in an array,
  - compact evenly-spaced bits in a word, and
  all in $O(1)$ machine word operations!

- Using these techniques, we can modify a B-tree to work strictly faster than a conventional BST, provided that we store tiny keys.

- Using these techniques, can we compute the most-significant bit of a machine word in time $O(1)$, independent of the machine word size.

- And all of this flows from one source: word-level parallelism inside of the processor!
What’s Next

• Can we build a data structure for integers that is strictly better than a binary search tree?
• The answer is yes, and it’s called a fusion tree.
• Today’s exploration provides the techniques we’ll use to build the fusion tree. We just need a few more insights to get us there!
Next Time

- **Patricia Codes**
  - Compressing a small number of big integers into a small number of small integers.

- **Fusion Trees**
  - Combining all these techniques together!