CRaft: Building High-Performance Consensus Protocols with Accurate Clocks

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• CRaft: a multi-leader extension to Raft enabled by accurate clocks



#### State Machines

- Maintain internal states
- Respond to external requests
- Examples: databases, storage systems



• How do we make them reliable?

### **Replicated State Machines**



- Consensus: ensures all servers agree on the same log
- Continues to operate if at least a majority of servers are up

### The Raft Consensus Protocol

- A widely used consensus protocol
- Leader-based
- Benefits: simple and efficient
- Limitation: leader is the bottleneck for throughput and scalability



Diego Ongaro and John Ousterhout. 2014. In search of an understandable consensus algorithm. In USENIX Annual Technical Conference. 305–319.

#### Limitations with Single Leader

• Single leader limits throughput and scalability



## Challenge in a Multi-Leader Protocol



- Challenge: how to coordinate leaders?
- Solution: agreement on time => agreement on order

#### **Clock Synchronization**

- Achieving agreement on time is not trivial in a distributed system
- Huygens: a software clock synchronization system

Distribution of clock offsets between servers (20 machines on CloudLab)

Percentile	90th	99th	99.9th	max
Clock offset	7us	11us	15us	26us



Yilong Geng, Shiyu Liu, Zi Yin, Ashish Naik, Balaji Prabhakar, Mendel Rosenblum, and Amin Vahdat. Exploiting a natural network effect for scalable, fine-grained clock synchronization. In NSDI 2018. 81–94.

### Our Approach: CRaft



### The CRaft Consensus Protocol

#### **CRaft Overview**





#### **Timestamp Management**



- CRaft guarantees monotonically increasing timestamps in each log
- Safe time: indicates how up-to-date a log is

#### Safe Times



# Merging



- Merge up to the smallest safe time
- CRaft ensures merged log in monotonically increasing timestamp order



• Fast path: respond to clients early for certain write operations

### Evaluation

### **Experiment Setup**

- Implementation
  - Based on HashiCorp Raft a popular and well-optimized implementation
- Environment
  - CloudLab, single data center
- Workload
  - In-memory key-value store
  - Multiple clients send get or set requests concurrently

### Throughput vs Cluster Size



• Up to ~2x read and ~2.5x write throughput compared to Raft

### Latency vs Throughput

Average latency vs throughput (3 servers)

35 35 Raft Raft 99th Percentile Latency (ms) CRaft - read CRaft - read . 30 CRaft - write CRaft - write 25 25 20 Performance gain under high load 15 10 0 5 5 Load increases 0 -0 20 60 80 100 120 100 0 20 40 60 80 120 40 Load increases Throughput (k ops/s) Throughput (k ops/s)

• CRaft improves throughput and latency under high load

99th percentile latency vs throughput (3 servers)

### Performance vs Number of Clients



• NTP precision: ~20ms, Huygens: ~20us

#### Conclusion



• Accurate clocks enable better performance and/or consistency

# Thank you!