EE270

Large scale matrix computation, optimization and learning

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Randomized Linear Algebra and Optimization Lecture 10: Leverage Scores and Basic Inequality Method

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Projected Least Squares Problems

Left-sketching

Form SA and Sb where $S \in \mathbb{R}^{m \times n}$ is a random projection matrix

Solve the smaller problem

$$\min_{x\in\mathbb{R}^d}\|SAx-Sb\|_2^2$$

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using any classical method.
 Direct method complexity md²

- We minimize $\tilde{x} = \arg \min \|S(Ax b)\|_2^2$
- x_{LS} minimizes $||Ax b||_2^2$
- How far is \tilde{x} from x_{LS} ?
- Step 1. Establish two optimality (in)equalities for these variables
- $||Ax_{LS} b||_2^2 \le ||Ax' b||_2^2$ for any x', i.e., $A^T(Ax_{LS} b) = 0$

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• $||S(A\tilde{x}-b)||_2^2 \le ||S(Ax_{LS}-b)||_2^2$

- We minimize $\tilde{x} = \arg \min \|S(Ax b)\|_2^2$
- x_{LS} minimizes $||Ax b||_2^2$
- How far is x from x_{LS}?
- Step 1. Establish two optimality (in)equalities for these variables
- $||Ax_{LS} b||_2^2 \le ||Ax' b||_2^2$ for any x', i.e., $A^T(Ax_{LS} b) = 0$
- $||S(A\tilde{x}-b)||_2^2 \le ||S(Ax_{LS}-b)||_2^2$
- Step 2. Define error Δ = x̃ x_{LS} and re-write these inequalities in terms of δ

$$||SA\Delta||_2^2 \le 2b^{\perp T}(S^TS - I)A\Delta$$

Step 3. Argue $S^T S \approx I$

Basic Inequality Method $\|SA\Delta\|_2^2 \le 2b^{\perp T}(S^T S - I)A\Delta$ $\max_{\Delta} \left| \frac{\|SA\Delta\|_2^2}{\|A\Delta\|_2^2} - 1 \right| = \max_{z} \left| \frac{\|SUz\|_2^2}{\|z\|_2^2} - 1 \right|$ $= \max_{z} \left| \frac{z^T}{\|z\|_2} (U^T S^T S U - I) \frac{z}{\|z\|_2} \right|$ $= \sigma_{zev}^2 (U^T S^T S U - I)$

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Basic Inequality Method $||SA\Delta||_2^2 \le 2b^{\perp T}(S^T S - I)A\Delta$ $\max_{\Delta} \left| \frac{||SA\Delta||_2^2}{||A\Delta||_2^2} - 1 \right| = \max_{z} \left| \frac{||SUz||_2^2}{||z||_2} - 1 \right|$ $= \max_{z} \left| \frac{z^T}{||z||_2} (U^T S^T S U - I) \frac{z}{||z||_2} \right|$ $= \sigma_{z}^2 \dots (U^T S^T S U - I)$

Approximate matrix multiplication (AMM):

$$\sigma_{\max}(U^T S^T S U - I) \le \|U^T S^T S U - \underbrace{U^T U}_{I}\|_F \le \epsilon \|U^T U\|_F^2$$

► This is called a Subspace Embedding we can rescale \(\epsilon\) to get \(\sigma_{max}(U^TS^TSU - I)) \le \(\epsilon\) for appropriate \(m\)

$$||SA\Delta||_2^2 \le 2b^{\perp T}(S^TS - I)A\Delta$$

Now consider the right handside

$$= 2b^{\perp T}(S^{T}S - I)UU^{T}A\Delta \leq 2\|b^{\perp T}(S^{T}S - I)UU^{T}\|_{2}\|A\Delta\|$$

► AMM again: $\|b^{\perp}S^{\top}SUU^{\top} - b^{\perp}UU^{\top}\|_{F} \le \frac{\epsilon}{\sqrt{m}}\|b^{\perp}\|_{F}\|UU^{\top}\|_{F} \le \frac{\epsilon}{\sqrt{m}}f(x_{LS})\sqrt{d}$

$$||SA\Delta||_2^2 \le 2b^{\perp T}(S^TS - I)A\Delta$$

Now consider the right handside

$$= 2b^{\perp T}(S^{T}S - I)UU^{T}A\Delta \leq 2\|b^{\perp T}(S^{T}S - I)UU^{T}\|_{2}\|A\Delta\|$$

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•
$$||SA\Delta||_2^2 \le 2b^{\perp T}(S^TS - I)A\Delta$$

• Summarizing two bounds:

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$$\|SA\Delta\|_{2}^{2} \leq 2b^{\perp T}(S^{T}S - I)A\Delta$$
Summarizing two bounds:
$$(1) \max_{\Delta} \left| \frac{\|SA\Delta\|_{2}^{2}}{\|A\Delta\|_{2}^{2}} - 1 \right| \leq \epsilon'$$

$$(2) 2b^{\perp T}(S^{T}S - I)A\Delta \leq \frac{\epsilon}{\sqrt{m}}f(x_{LS})\sqrt{d}\|A\Delta\|_{2}$$

$$(1) \text{ implies } -\epsilon'\|A\Delta\|_{2}^{2} \leq \|SA\Delta\|_{2}^{2} - \|A\Delta\|_{2}^{2} \leq \epsilon'\|A\Delta\|_{2}^{2}$$

$$\text{hence } (1 - \epsilon')\|A\Delta\|_{2}^{2} \leq \|SA\Delta\|_{2}^{2}$$

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$$(1) \text{ implies } -\epsilon'\|A\Delta\|_{2}^{2} \leq \|SA\Delta\|_{2}^{2} - \|A\Delta\|_{2}^{2} \leq \epsilon'\|A\Delta\|_{2}^{2}$$
hence $(1 - \epsilon')\|A\Delta\|_{2}^{2} \leq \|SA\Delta\|_{2}^{2}$
Plugging in: $(1 - \epsilon')\|A\Delta\|_{2}^{2} \leq \frac{\epsilon}{\sqrt{m}}f(x_{LS})\sqrt{d}\|A\Delta\|_{2}$

$$\|A\Delta\|_{2} \leq \frac{\epsilon}{1-\epsilon'}f(x_{LS})\frac{\sqrt{d}}{\sqrt{m}}$$

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Leverage Scores

► Intuition: Approximate Matrix Multiplication for $U^T U$ i.e, $\|U^T S^T S U - U^T U\|_F = \|U^T S^T S U - I\|_F \le \epsilon$

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implies Least Squares cost approximation

Leverage Scores

► Intuition: Approximate Matrix Multiplication for $U^T U$ i.e, $\|U^T S^T S U - U^T U\|_F = \|U^T S^T S U - I\|_F \le \epsilon$

implies Least Squares cost approximation

- We can pick a sampling matrix S
- Importance sampling: proportional to the rows norms of U
- Leverage scores: $\ell_i := ||u_i||_2^2$ for i = 1, ..., n
- $\sum_i \ell_i = \sum_i ||u_i||_2^2 = ||U||_F^2 = trU^T U = trI_d = d$ when A is full column rank

- Sampling probabilities: $p_i = \frac{1}{d} ||u_i||_2^2$ $\sum_i p_i = 1$
- Can be non-uniform or uniform A = [1; 0]

Fast Johnson Lindenstrauss Transform

- Let *H* be the $n \times n$ Hadamard matrix
- Generate an n × n diagonal matrix of random ±1 uniform signs

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- Uniform $m \times n$ sub-sampling matrix P scaled with $\frac{\sqrt{n}}{\sqrt{m}}$
- Let S = PHD.
- Note that $\mathbb{E}S^T S = I$

FJLT Preconditions Leverage Scores

Fix a set X of n vectors in d-dimension. With high probability $\max_{x \in X} \|HDX\|_{\infty} \leq \sqrt{\frac{\log(n)}{d}}$

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FJLT Preconditions Leverage Scores

 Fix a set X of n vectors in d-dimension. With high probability max_{x∈X} ||HDX||_∞ ≤ √ ^{log(n)}/_d Apply HD to data A
 PHDA is uniformly sampled HDA Leverage scores of HDU are near uniform uniform sampling works!

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Questions?