

EE270

Large scale matrix computation, optimization and learning

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Randomized Linear Algebra and Optimization

Lecture 17: Randomized Singular Value Decomposition and CX Decomposition

Randomized Matrix Decompositions

- ▶ Suppose that A is an $n \times d$ data matrix of rank r
- ▶ Singular Value Decomposition (SVD) provides the best rank k approximation:

Let $A = U\Sigma V^T = \sum_{i=1}^r \sigma_i u_i v_i^T$ where $\{\sigma_i\}_{i=1}^r$ are the singular values sorted in non-increasing order

Define $A_k := U_k \Sigma_k V_k^T := \sum_{i=1}^k \sigma_i u_i v_i^T$. We have

$$\|A - A_k\|_2 \leq \sigma_{k+1}$$

- ▶ computational cost of computing the SVD is $O(nd^2)$ for $n \geq d$

Randomized Matrix Decompositions

- ▶ Given a large matrix $A \in \mathbb{R}^{n \times d}$
- ▶ idea: sample some columns of A to get $C \in \mathbb{R}^{n \times m}$

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- ▶ form an approximation of A using these sampled columns
 $C = AS$

$$\min_X \|CX - A\|_F^2 = \min_X \|ASX - A\|_F^2$$

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$$\min_X \|CX - A\|_F^2 = \min_X \|ASX - A\|_F^2$$

- ▶ column-wise decomposable problem

$$\arg \min_{X^{(j)}} \sum_{k=1}^d \|ASX^{(k)} - A^{(k)}\|_2^2 = (AS)^\dagger A^{(k)}$$

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$$\arg \min_X \|ASX - A\|_F^2 = (AS)^\dagger A$$

- ▶ matrix A is approximated by $(AS)(AS)^\dagger A = CC^\dagger A$

Randomized Matrix Decompositions: Spectral Norm Error

- ▶ Given a large matrix $A \in \mathbb{R}^{n \times d}$
- ▶ idea: sample some columns of A to get $C \in \mathbb{R}^{n \times m}$
- ▶ approximation $CC^\dagger A \approx A$
- ▶ calculate top left singular values of $C \approx U_k \Sigma_k V_k^T$
- ▶ then we have $CC^\dagger \approx U_k U_k^T$

Lemma 1 (Drineas et al. 2006)

$$\|A - U_k U_k^T A\|_2^2 \leq \|A - U_A U_A^T A\|_2^2 + 2\|AA^T - CC^T\|_2$$

- ▶ first term is the approximation error of the exact SVD
- ▶ second term is the spectral norm approximate matrix multiplication error
- ▶ approximate matrix multiplication results can be used

Randomized Matrix Decompositions: Spectral Norm Error

- ▶ $C = AS$, approximation $CC^\dagger A \approx A$
- ▶ calculate top left singular values of $C \approx U_k \Sigma_k V_k^T$
- ▶ approximate $A \approx U_k U_k^T A$

Proof of Lemma 1

$$\begin{aligned} & \|A - U_k U_k^T A\|_2^2 \\ &= \max_{\|x\|_2=1} \|x^T (A - U_k U_k^T A)\|_2 \\ &= \max_{\substack{\|y\|_2=\|z\|_2=1, y \in U_k, z \in U_k^\perp \\ \alpha^2 + \beta^2 = 1}} \|(\alpha y + \beta z)^T (A - U_k U_k^T A)\|_2 \\ &\leq \max_{\|z\|_2=1, z \in U_k^\perp} \|z^T (A - U_k U_k^T A)\|_2 + \\ &\quad \max_{\|y\|_2=1, y \in U_k} \|y^T (A - U_k U_k^T A)\|_2 \\ &= \max_{\|z\|_2=1, z \in U_k^\perp} \|z^T (A - U_k U_k^T A)\|_2 \end{aligned}$$

Randomized Matrix Decompositions: Spectral Norm Error

Proof of Lemma 1 cont'd

taking squares

$$\begin{aligned} & \max_{\|z\|_2=1, z \in U_k^\perp} \|z^T A\|_2^2 \\ &= \max_{\|z\|_2=1, z \in U_k^\perp} z^T C C^T z + z^T (A A^T - C C^T) z \\ &\leq \max_{\|z\|_2=1, z \in U_k^\perp} \sigma_{k+1}^2(C) + \|A A^T - C C^T\|_2 \\ &\leq \max_{\|z\|_2=1, z \in U_k^\perp} \sigma_{k+1}^2(A) + 2\|A A^T - C C^T\|_2 \end{aligned}$$

Randomized Matrix Decompositions: Spectral Norm Error

Proof of Lemma 1 cont'd

taking squares

$$\begin{aligned} & \max_{\|z\|_2=1, z \in U_k^\perp} \|z^T A\|_2^2 \\ &= \max_{\|z\|_2=1, z \in U_k^\perp} z^T C C^T z + z^T (A A^T - C C^T) z \\ &\leq \max_{\|z\|_2=1, z \in U_k^\perp} \sigma_{k+1}^2(C) + \|A A^T - C C^T\|_2 \\ &\leq \max_{\|z\|_2=1, z \in U_k^\perp} \sigma_{k+1}^2(A) + 2\|A A^T - C C^T\|_2 \end{aligned}$$

- ▶ where we used a matrix perturbation result
 $\sigma_{k+1}(C C^T) - \sigma_{k+1}(A A^T) \leq \|A A^T - C C^T\|_2$
- ▶ Hoffman–Wielandt inequality:
- ▶ $\max_k |\sigma_k(Q + E) - \sigma_k(Q)| \leq \|E\|_2$

Randomized Matrix Decompositions: Frobenius Norm Error

- ▶ Given a large matrix $A \in \mathbb{R}^{n \times d}$
- ▶ idea: sample some columns of A to get $C \in \mathbb{R}^{n \times m}$
- ▶ approximation $CC^\dagger A \approx A$
- ▶ calculate top left singular values of $C = U_k \Sigma_k V_k^T$

Lemma 2 (Drineas et al. 2006)

$$\|A - U_k U_k^T A\|_F^2 \leq \|A - U_A U_A^T A\|_F^2 + 2\sqrt{k} \|AA^T - CC^T\|_F$$

- ▶ approximate matrix multiplication results can be used

Randomized Singular Value Decomposition

- ▶ Given a large matrix $A \in \mathbb{R}^{n \times d}$
- ▶ idea: sample some columns of A to get $C \in \mathbb{R}^{n \times m}$
- ▶ $C = AS$
- ▶ approximation $CC^\dagger = (AS)(AS)^\dagger A \approx A$
- ▶ calculate QR decomposition of $AS = QR$
- ▶ then $QQ^T A \approx A$, i.e., Q approximate the range space of A
- ▶ calculate the SVD $Q^T A = U\Sigma V^T$
- ▶ approximate SVD of A is $A \approx (QU)\Sigma V^T$

Randomized Singular Value Decomposition

- ▶ Given a large matrix $A \in \mathbb{R}^{n \times d}$
- ▶ idea: sample/sketch some columns of A to get $C \in \mathbb{R}^{n \times m}$
- ▶ $C = AS$
- ▶ approximation $CC^\dagger = (SA)(SA)^\dagger A \approx A$
- ▶ calculate QR decomposition of $SA = QR$
- ▶ then $QQ^T A \approx A$, i.e., Q approximate the range space of A
- ▶ calculate the SVD $Q^T A = U\Sigma V^T$
- ▶ approximate SVD of A is $A \approx (QU)\Sigma V^T$

Randomized Singular Value Decomposition: Analysis

- ▶ Given a large matrix $A \in \mathbb{R}^{n \times d}$
- ▶ Generate a Gaussian sketching matrix $S \in d \times m$
- ▶ $C = AS$
- ▶ approximation $CC^\dagger = (AS)(AS)^\dagger A \approx A$
- ▶ calculate QR decomposition of $C = AS = QR$
- ▶ then $QQ^T A \approx A$, i.e., Q approximate the range space of A
- ▶ calculate the SVD $Q^T A = U\Sigma V^T$
- ▶ approximate SVD of A is $A \approx (QU)\Sigma V^T$
- ▶ **Lemma** (Halko et al. 2009)

$$\mathbb{E}\|A - QQ^T A\|_2 \leq \left(1 + \frac{4\sqrt{m}}{m - k - 1} \sqrt{\min(n, d)}\right) \sigma_{k+1}$$

Randomized Singular Value Decomposition: Comparison

- ▶ Generate a Gaussian sketching matrix $S \in d \times m$
- ▶ calculate QR decomposition of $C = AS = QR$
- ▶ calculate the SVD $Q^T A = U \Sigma V^T$
- ▶ approximate SVD of A is $A \approx (QU) \Sigma V^T$
- ▶ **Lemma** (Halko et al. 2009)

$$\mathbb{E} \|A - QQ^T A\|_2 \leq \left(1 + \frac{4\sqrt{m}}{m - k - 1} \sqrt{\min(n, d)} \right) \sigma_{k+1}$$

- ▶ Exact SVD of $A = U_A \Sigma_A V_A^T$ yields

$$\|A - U_A^k (U_A^k)^T A\|_2 \leq \sigma_{k+1}$$

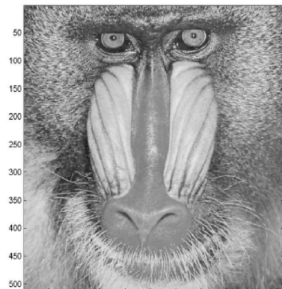
Low-rank matrix approximations

- ▶ Singular Value Decomposition (SVD)
- ▶ $A = U\Sigma V^T$
- ▶ takes $O(nd^2)$ time for $A \in R^{n \times d}$
- ▶ best rank- k approximation is $A_k := U_k \Sigma_k V_k^T = \sum_{i=1}^k \sigma_i u_i v_i^T$
- ▶ $\|A - A_k\|_2 \leq \sigma_{k+1}$

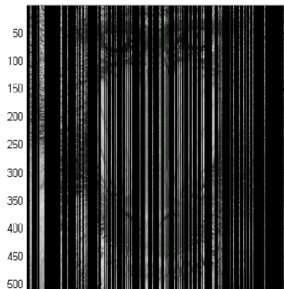
Randomized low-rank matrix approximations

- ▶ Randomized (SVD)
- ▶ approximation C (e.g. a subset of the columns of A)
- ▶ $AA^T \approx CC^T$
- ▶ $\tilde{A}_k = CC^\dagger A$ is a randomized rank- k approximation
- ▶ $\|A - \tilde{A}_k\|_2^2 \leq \sigma_{k+1}^2 + \epsilon \|A\|_2^2$

Randomized low-rank approximation example

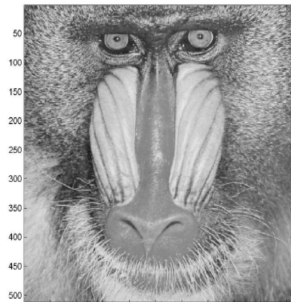


original data

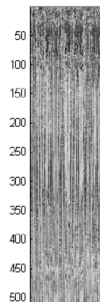


column subsampled data

Randomized low-rank approximation example

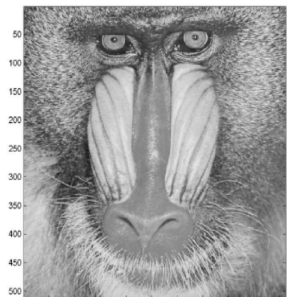


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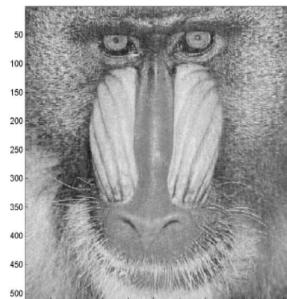


sketch

Randomized low-rank approximation example



original data



randomized low rank approximation