

EE270

Large scale matrix computation, optimization and learning

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Randomized Linear Algebra and Optimization
Lecture 18: Generalized Least Squares Problems,
Randomized Low Rank Approximations and Power
Iteration

Recap: Low-rank matrix approximations

- ▶ Singular Value Decomposition (SVD)

$$A = U\Sigma V^T$$

takes $O(nd^2)$ time for $A \in R^{n \times d}$

- ▶ best rank- k approximation $A_k := U_k \Sigma_k V_k^T = \sum_{i=1}^k \sigma_i u_i v_i^T$

$$\|A - A_k\|_2 \leq \sigma_{k+1}$$

Recap: Randomized low-rank matrix approximations

idea: sample some rows/sketch $A \in \mathbb{R}^{n \times d}$ to get $C \in \mathbb{R}^{n \times m}$
 $C = AS$ where $S \in \mathbb{R}^{d \times m}$ is a sampling/sketching matrix

- ▶ we have an approximate matrix multiplication $AA^T \approx CC^T$.
Next, consider the best approximation of A
in the range of $C = AS$

$$\min_X \|CX - A\|_F$$

also called CX decomposition

- ▶ $\tilde{A}_m := CX^* = CC^\dagger A$ is a randomized rank- m approximation

$$(AS)(AS)^\dagger A \approx A$$

Recap: Randomized Singular Value Decomposition

- ▶ CX decomposition provides the approximation

$$(AS)(AS)^\dagger A \approx A$$

- ▶ calculate QR decomposition of $AS = QR$
 $QQ^T A \approx A$, i.e., Q approximates the range space of A
- ▶ calculate the SVD $Q^T A = U\Sigma V^T$
- ▶ approximate SVD of A is $A \approx \underbrace{(QU)}_{U'} \Sigma V^T$

note that $U' := QU$ and V are orthogonal and Σ is diagonal
we have an approximation of the left and right singular vectors and singular values of A

Analysis of Randomized Low Rank Approximations

- ▶ CX decomposition: form sketch AS , and find the best approximation of A in the range of AS

$$X^* = \arg \min_X \|ASX - A\|_F^2 = (AS)^\dagger A$$

- ▶ approximation $ASX^* = (AS)(AS)^\dagger A \approx A$
yields randomized SVD : $AS = QR$ and $Q^T A = U\Sigma V^T$
- ▶ Let $A = U\Sigma V^T$ and $A_k = \sum_{i=1}^k \sigma_k u_k v_k^T$, i.e.,
best rank- k approximation of A
note that

$$\begin{aligned} \|AS \underbrace{(AS)^\dagger A}_{X^*} - A\|_F^2 &\leq \|AS(A_k S)^\dagger A_k - A\|_F^2 \\ &= \|A_k^T (S^T A_k^T)^\dagger S^T A^T - A^T\|_F^2 \end{aligned}$$

Analysis of Randomized Low Rank Approximations

- ▶ approximation error

$$\begin{aligned}\|AS \underbrace{(AS)^\dagger A}_{X^*} - A\|_F^2 &\leq \|AS(A_k S)^\dagger A_k - A\|_F^2 \\ &= \|A_k^T (S^T A_k^T)^\dagger S^T A^T - A^T\|_F^2 \\ &= \|A_k^T \tilde{X} - A^T\|_F^2\end{aligned}$$

where

$$\tilde{X} := \arg \min_X \|S^T A_k^T X - S^T A^T\|_F^2$$

Analysis of Randomized Low Rank Approximations

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$$\begin{aligned}\|AS \underbrace{(AS)^\dagger}_X A - A\|_F^2 &\leq \|AS(A_k S)^\dagger A_k - A\|_F^2 \\ &= \|A_k^T (S^T A_k^T)^\dagger S^T A^T - A^T\|_F^2 \\ &= \|A_k^T \tilde{X} - A^T\|_F^2\end{aligned}$$

where

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- ▶ identical to left sketching the Generalized Least Squares problem

$$\min_X \|A_k^T X - A^T\|_F^2$$

Generalized Least Squares Problems

$$\min_X \|AX - B\|_F^2$$

- ▶ Least Squares problem with multiple right-hand-sides

$$B = [b_1, \dots, b_r]$$

$$X = [x_1, \dots, x_r]$$

$$\min_{x_1, \dots, x_r} \sum_{i=1}^r \|Ax_i - b_i\|_2^2$$

- ▶ optimal solution

$$\begin{aligned} X^* &= [x_1^*, \dots, x_r^*] \\ &= [A^\dagger b_1, \dots, A^\dagger b_r] \\ &= A^\dagger B \end{aligned}$$

Left Sketching Generalized Least Squares Problems

- ▶ original problem

$$X^* := \arg \min_X \|AX - B\|_F^2$$

- ▶ form sketches of the data SA and SB , e.g., uniform row sampling, weighted sampling, Gaussian, ± 1 i.i.d, CountSketch, FJLT...

$$\hat{X} := \arg \min_X \|SAX - SB\|_F^2$$

$$\begin{aligned}\hat{X}_i &= \arg \min_{x_i} \|SAx_i - Sb_i\|_2^2 \\ &= (SA)^\dagger(Sb_i)\end{aligned}$$

- ▶ left-sketch applied to simple Least Squares problem $\min_{x_i} \|Ax_i - b_i\|_2^2$

Recall Gaussian Sketch Analysis

- ▶ Let $A \in \mathbb{R}^{n \times d}$, $S \in \mathbb{R}^{m \times n}$ be i.i.d. Gaussian

$$x^* := \arg \min_{x \in \mathbb{R}^d} \underbrace{\|Ax - b\|_2^2}_{f(x)} \quad \text{and} \quad \tilde{x} = \arg \min_{x \in \mathbb{R}^d} \|SAx - Sb\|_2^2$$

- ▶ Conditioned on the matrix SA

$$A(\tilde{x} - x^*) \sim N\left(0, \frac{f(x^*)}{m} A(A^T S^T SA)^{-1} A\right)$$

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- ▶ Conditioned on the matrix SA

$$A(\tilde{x} - x^*) \sim N\left(0, \frac{f(x^*)}{m} A(A^T S^T SA)^{-1} A\right)$$

- ▶ taking expectation over SA , and using $\mathbb{E}[(A^T S^T SA)^{-1}] = (A^T A)^{-1} \frac{m}{m-d-1}$ we get

$$\begin{aligned} \mathbb{E} \|A(\tilde{x} - x^*)\|_2^2 &= \frac{f(x^*)}{m-d-1} \text{tr} A(A^T A)^{-1} A \\ &= f(x^*) \frac{\text{rank}(A)}{m-d-1} = f(x^*) \frac{d}{m-d-1} \end{aligned}$$

Left Sketching Generalized Least Squares Problems

- ▶ original problem and left-sketch

$$X^* := \arg \min_X \|AX - B\|_F^2 \quad \text{and} \quad \hat{X} := \arg \min_X \|SAX - SB\|_F^2$$

- ▶ x_i : i -th column of \hat{X} satisfies

$$\hat{x}_i = \arg \min_{x_i} \|SAx_i - Sb_i\|_2^2$$

- ▶ For a Gaussian sketching matrix S we have

$$\mathbb{E} \|A(\hat{x}_i - x_i^*)\|_2^2 = \|Ax_i^* - b_i\|_2^2 \frac{d}{m - d - 1}$$

implies

$$\begin{aligned} \mathbb{E} \|A(\hat{X} - X^*)\|_F^2 &= \sum_{i=1}^r \|Ax_i^* - b_i\|_2^2 \frac{d}{m - d - 1} \\ &= \|AX^* - B\|_F^2 \frac{d}{m - d - 1} \end{aligned}$$

Left Sketching Optimality Gap

- ▶ suppose that $\mathbf{rank}(A) = r$
original problem and left-sketch

$$X^* := \arg \min_X \|AX - B\|_F^2 \quad \text{and} \quad \hat{X} := \arg \min_X \|SAX - SB\|_F^2$$

$$\mathbb{E} \|A(\hat{X} - X^*)\|_F^2 = \|AX^* - B\|_F^2 \frac{r}{m - r - 1}$$

$$\begin{aligned} \mathbb{E} \|A\hat{X} - B\|_F^2 &= \mathbb{E} \|AX^* - B + A(\hat{X} - X^*)\|_F^2 \\ &= \|AX^* - B\|_F^2 + \mathbb{E} \|A(\hat{X} - X^*)\|_F^2 \\ &= \|AX^* - B\|_F^2 \left(1 + \frac{r}{m - r - 1}\right) \\ &= \|AX^* - B\|_F^2 \frac{m - 1}{m - r - 1} \end{aligned}$$

Back to Randomized Low Rank Approximations

- ▶ approximation error

$$\begin{aligned}\mathbb{E}\|AS \underbrace{(AS)^\dagger A}_{X^*} - A\|_F^2 &\leq \mathbb{E}\|AS(A_k S)^\dagger A_k - A\|_F^2 \\ &= \|A_k^T (S^T A_k^T)^\dagger S^T A^T - A^T\|_F^2 \\ &= \mathbb{E}\|A_k^T \tilde{X} - A^T\|_F^2 \\ &\leq \frac{m-1}{m-k-1} \|A_k^T (A_k^T)^\dagger A^T - A^T\|_F^2 \\ &\leq \frac{m-1}{m-k-1} \|A(A_k A_k^\dagger - I)\|_F^2 \\ &\leq \frac{m-1}{m-k-1} \|A_k - A\|_F^2\end{aligned}$$

Randomized Low Rank Approximation and Randomized SVD Error Bound

- ▶ CX decomposition and randomized SVD

$$AS(AS)^\dagger A \approx A$$

- ▶ final Frobenious norm error bound

$$\mathbb{E} \|AS(AS)^\dagger A - A\|_F^2 \leq \frac{m-1}{m-k-1} \|A_k - A\|_F^2$$

valid for any $k \in \{1, \dots, \mathbf{rank}(A)\}$

Randomized Low Rank Approximation and Randomized SVD Error Bound

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$$\mathbb{E} \|AS(AS)^\dagger A - A\|_F^2 \leq \frac{m-1}{m-k-1} \|A_k - A\|_F^2$$

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- ▶ define the oversampling factor $\ell := m - k - 1$

$$\|AS(AS)^\dagger A - A\|_F^2 \leq \left(1 + \frac{k}{\ell}\right) \|A_k - A\|_F^2$$

Reducing the Error: Power Iteration

- ▶ error bounds depend on tail singular values

$$\|A_k - A\|_F^2 = \sum_{j=k+1}^{\text{rank}(A)} \sigma_j^2$$

- ▶ idea: compute the sketch of $(AA^T)^q A$

$$C = (AA^T)^q A S$$

where q is an integer parameter

$$CC^\dagger A \approx A$$

CC^\dagger approximates the range of A better for $q \geq 1$

- ▶ singular values of $(AA^T)^q A$ are $\sigma_i(A)^{2q+1}$
where $\sigma_i(A)$ are the singular values of A

Connection to the Classical Power Iteration

- ▶ suppose that $A^T A$ has a unique maximal eigenvector
- ▶ Starting from a nonzero vector x_0 , the iterations

$$x_{t+1} = \frac{A^T A x_t}{\|A^T A x_t\|_2}$$

converges to a multiple of the maximal eigenvector

Simultaneous (QR) Iteration

- ▶ suppose that the top- m eigenvalues of $A^T A$ are distinct
- ▶ Starting from a matrix X_0 of rank m , e.g., a random $d \times m$ matrix, the iterations

$$Q_t R_t = X_t$$

$$X_{t+1} = A^T A Q_t$$

where $Q_t R_t$ is the QR factorization of X_t ,

converges to an orthonormal basis for the top- m eigenspace of $A^T A$

- ▶ note that the first iteration computes the QR decomposition $A^T A S = QR$ where S is random. This is similar to the CR-decomposition and randomized SVD approach, where the QR decomposition of AS serves as a crude approximation of the top eigenspace basis