

due May 10, Friday 23:59pm

**EE364b Homework 5**

4.0 *Maximum volume ellipsoid vs Chebyshev center method.* Consider the convex set

$$\mathcal{C} = \{x \mid Ax \preceq b\},$$

where  $A \in \mathbf{R}^{n \times d}$  and  $b \in \mathbf{R}^d$ . The data files `Amatrix` and `bvector` are available on Canvas.

- (a) Find the center of the maximum volume ellipsoid in  $\mathcal{C}$  and the center of the largest Euclidean ball in  $\mathcal{C}$ . You may use CVX/CVXPY. *Hint: See 364a slides for calculating the maximum volume ellipsoid.*
- (b) Denote the two centers (vectors in  $\mathbf{R}^d$ ) in part (a) by  $x_{\text{ellipsoid}}$  and  $x_{\text{ball}}$  respectively. Let  $g \in \mathbf{R}^d$  be the all-ones vector. We will consider the cuts  $g^T(x - x_{\text{ball}}) \geq 0$  and  $g^T(x - x_{\text{ellipsoid}}) \geq 0$ . Estimate the volume ratios

$$R_{\text{ellipsoid}} := \frac{\text{vol}(\{g^T(x - x_{\text{ellipsoid}}) \geq 0\} \cap \mathcal{C})}{\text{vol}(\mathcal{C})},$$

and

$$R_{\text{ball}} := \frac{\text{vol}(\{g^T(x - x_{\text{ball}}) \geq 0\} \cap \mathcal{C})}{\text{vol}(\mathcal{C})},$$

by generating  $M = 10^6$  i.i.d. uniformly distributed random vectors in  $[-0.5, +0.5]^d$  (i.e.,  $x = \text{rand}(d, 1) - 0.5$  for  $M$  trials). *Hint:* Let  $M_{\mathcal{C}}$  be number of random vectors that satisfy  $Ax \preceq b$ . Let  $M_{\text{ellipsoid}}$  be the number of random vectors that satisfy  $Ax \preceq b$  and  $g^T(x - x_{\text{ellipsoid}}) \geq 0$ . Similarly, let  $M_{\text{ball}}$  be the number of random vectors that satisfy  $Ax \preceq b$  and  $g^T(x - x_{\text{ball}}) \geq 0$ . The volume ratios can be estimated by

$$R_{\text{ellipsoid}} \approx \frac{M_{\text{ellipsoid}}}{M_{\mathcal{C}}},$$

and

$$R_{\text{ball}} \approx \frac{M_{\text{ball}}}{M_{\mathcal{C}}}.$$

4.1 *Kelley's cutting-plane algorithm.* We consider the problem of minimizing a convex function  $f : \mathbf{R}^n \rightarrow \mathbf{R}$  over some convex set  $C$ , assuming we can evaluate  $f(x)$  and find a subgradient  $g \in \partial f(x)$  for any  $x$ . Suppose we have evaluated the function and a subgradient at  $x^{(1)}, \dots, x^{(k)}$ . We can form the piecewise-linear approximation

$$\hat{f}^{(k)}(x) = \max_{i=1, \dots, k} (f(x^{(i)}) + g^{(i)T}(x - x^{(i)})),$$

which satisfies  $\hat{f}^{(k)}(x) \leq f(x)$  for all  $x$ . It follows that

$$L^{(k)} = \inf_{x \in C} \hat{f}^{(k)}(x) \leq p^*,$$

where  $p^* = \inf_{x \in C} f(x)$ . Since  $\hat{f}^{(k+1)}(x) \geq \hat{f}^{(k)}(x)$  for all  $x$ , we have  $L^{(k+1)} \geq L^{(k)}$ .

In Kelley's cutting-plane algorithm, we set  $x^{(k+1)}$  to be any point that minimizes  $\hat{f}^{(k)}$  over  $x \in C$ . The algorithm can be terminated when  $U^{(k)} - L^{(k)} \leq \epsilon$ , where  $U^{(k)} = \min_{i=1, \dots, k} f(x^{(i)})$ .

Use Kelley's cutting-plane algorithm to minimize the piecewise-linear function

$$f(x) = \max_{i=1, \dots, m} (a_i^T x + b_i)$$

that we have used for other numerical examples, with  $C$  the unit cube, *i.e.*,  $C = \{x \mid \|x\|_\infty \leq 1\}$ . Generate the same data we used before using

```
n = 20; % number of variables
m = 100; % number of terms
randn('state', 1);
A = randn(m, n);
b = randn(m, 1);
```

You can start with  $x^{(1)} = 0$  and run the algorithm for 40 iterations. Plot  $f(x^{(k)})$ ,  $U^{(k)}$ ,  $L^{(k)}$  and the constant  $p^*$  (on the same plot) versus  $k$ .

Repeat for  $f(x) = \|x - c\|_2$ , where  $c$  is chosen from a uniform distribution over the unit cube  $C$ . (The solution to this problem is, of course,  $x^* = c$ .)

4.4 *Minimum volume ellipsoid covering a half-ellipsoid.* In this problem we derive the update formulas used in the ellipsoid method, *i.e.*, we will determine the minimum volume ellipsoid that contains the intersection of the ellipsoid

$$\mathcal{E} = \{x \in \mathbf{R}^n \mid (x - x_c)^T P^{-1} (x - x_c) \leq 1\}$$

and the halfspace

$$\mathcal{H} = \{x \mid g^T (x - x_c) \leq 0\}.$$

We'll assume that  $n > 1$ , since for  $n = 1$  the problem is easy.

- (a) We first consider a special case:  $\mathcal{E}$  is the unit ball centered at the origin ( $P = I$ ,  $x_c = 0$ ), and  $g = -e_1$  ( $e_1$  is the first unit vector), so  $\mathcal{E} \cap \mathcal{H} = \{x \mid x^T x \leq 1, x_1 \geq 0\}$ . Let

$$\tilde{\mathcal{E}} = \{x \mid (x - \tilde{x}_c)^T \tilde{P}^{-1} (x - \tilde{x}_c) \leq 1\}$$

denote the minimum volume ellipsoid containing  $\mathcal{E} \cap \mathcal{H}$ . Since  $\mathcal{E} \cap \mathcal{H}$  is symmetric about the line through first unit vector  $e_1$ , it is clear (and not too hard to show) that  $\tilde{\mathcal{E}}$  will have the same symmetry. This means that the matrix  $\tilde{P}$  is diagonal, of the form  $\tilde{P} = \mathbf{diag}(\alpha, \beta, \beta, \dots, \beta)$ , and that  $\tilde{x}_c = \gamma e_1$  (where  $\alpha, \beta > 0$  and  $\gamma \geq 0$ ).

So now we have only three variables to determine:  $\alpha$ ,  $\beta$ , and  $\gamma$ . Express the volume of  $\tilde{\mathcal{E}}$  in terms of these variables, and also the constraint that  $\tilde{\mathcal{E}} \supseteq \mathcal{E} \cap \mathcal{H}$ . Then solve the optimization problem directly, to show that

$$\alpha = \frac{n^2}{(n+1)^2}, \quad \beta = \frac{n^2}{n^2-1}, \quad \gamma = \frac{1}{n+1}$$

(which agrees with the formulas we gave, for this special case).

*Hint.* To express  $\mathcal{E} \cap \mathcal{H} \subseteq \tilde{\mathcal{E}}$  in terms of the variables, it is necessary and sufficient for the conditions on  $\alpha$ ,  $\beta$ , and  $\gamma$  to hold on the boundary of  $\mathcal{E} \cap \mathcal{H}$ , *i.e.*, at the points

$$x_1 = 0, \quad x_2^2 + \dots + x_n^2 \leq 1,$$

or the points

$$x_1 \geq 0, \quad x_1^2 + x_2^2 + \dots + x_n^2 = 1.$$

- (b) Now consider the general case, stated at the beginning of this problem. Show how to reduce the general case to the special case solved in part (a).

*Hint.* Find an affine transformation that maps the original ellipsoid to the unit ball, and  $g$  to  $-e_1$ . Explain why minimizing the volume in these transformed coordinates also minimizes the volume in the original coordinates.

- (c) Finally, show that the volume of the ellipse  $\tilde{\mathcal{E}}$  satisfies  $\mathbf{vol}(\tilde{\mathcal{E}}) \leq e^{-\frac{1}{2n}} \mathbf{vol}(\mathcal{E})$ .

*Hint.* Compute the volume of the ellipse  $\mathcal{E}$  as a function of the eigenvalues of  $P$ , then use the results of parts (a) and (b) to argue that the volume computation can be reduced to the special case in part (a).