

# On Computing the Worst-Case Peak Gain of Linear Systems

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## Abstract

Based on the bounds due to Doyle and Boyd, we present simple upper and lower bounds for the  $\ell^1$ -norm of the ‘tail’ of the impulse response of finite-dimensional discrete-time linear time-invariant systems. Using these bounds, we may in turn compute the  $\ell^\infty$ -gain of these systems to any desired accuracy. By combining these bounds with results due to Khammash and Pearson, we derive upper and lower bounds for the worst-case  $\ell^\infty$ -gain of discrete-time systems with diagonal perturbations.

**Keywords:** SISO discrete-time LTI systems, computation of  $\ell^\infty$ -gain, discrete-time systems with diagonal perturbations, worst-case  $\ell^\infty$ -gain.

## 1 Notation

$\mathbf{Z}_+$ ,  $\mathbf{R}$ ,  $\mathbf{R}_+$  and  $\mathbf{C}$  denote the set of nonnegative integers, real numbers, nonnegative real numbers and complex numbers respectively. All the sequences in this note are defined over  $\mathbf{Z}_+$ . The  $\ell^\infty$ -norm of a complex-valued sequence  $u$  is defined as

$$\|u\|_\infty \triangleq \sup_{k \geq 0} |u(k)|.$$

Thus, the  $\ell^\infty$ -norm of a sequence is its peak value. The  $\ell^1$ -norm of a complex-valued sequence  $u$  is defined as

$$\|u\|_1 \triangleq \sum_{k \geq 0} |u(k)|.$$

For a matrix  $P \in \mathbf{R}^{n \times n}$ ,  $P^T$  stands for the transpose.  $\sigma_1(P), \sigma_2(P), \dots, \sigma_n(P)$  are the singular values of  $P$  in decreasing order.  $\rho(P)$  denotes the spectral radius, which is the maximum magnitude of the eigenvalues of  $P$ .  $I$  stands for the identity matrix, with size determined from context.

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## 2 Bounds for the $\ell^\infty$ -gain

Consider a stable, finite-dimensional discrete-time linear time-invariant (LTI) system described by the state equations

$$\begin{aligned} x(k+1) &= Ax(k) + bu(k), & x(0) &= 0, \\ y(k) &= cx(k) + du(k), \end{aligned} \tag{1}$$

where the input  $u(k) \in \mathbf{R}$ , the output  $y(k) \in \mathbf{R}$  and the state  $x(k) \in \mathbf{R}^n$ . We assume that  $\{A, b, c, d\}$  is minimal. The *impulse response* of system (1) is the real sequence given by

$$h(k) \triangleq \begin{cases} d, & k = 0, \\ cA^{k-1}b, & k > 0. \end{cases}$$

The  $\ell^\infty$ -gain of system (1), which is the largest possible peak value of the output  $y$  over all possible inputs  $u$  with a peak value of at most one, is just  $\|h\|_1$ :

$$\|h\|_1 = \sup_{\|u\|_\infty > 0} \frac{\|y\|_\infty}{\|u\|_\infty}.$$

$\|h\|_1$  is usually approximated by summing only a finite, typically large (say  $N$ ) number of terms:

$$S_N = \sum_{k=0}^N |h(k)| \leq \|h\|_1.$$

Obviously,  $S_N$  is a lower bound for  $\|h\|_1$ , and increases monotonically to  $\|h\|_1$  with increasing  $N$ . The ‘error’  $\|h\|_1 - S_N$  is just the  $\ell^1$  norm of the tail,  $\sum_{k>N} |h(k)|$ . Many simple bounds on this error are possible; for instance, if the poles of the system (1) are distinct, we may write down a residue expansion for the impulse response  $h(k)$ :

$$h(k) = \begin{cases} d, & k = 0, \\ \sum_{i=1}^n r_i p_i^{k-1}, & k > 0. \end{cases}$$

where  $p_1, p_2, \dots, p_n$  are the distinct poles of the system and  $r_i$  are the residues (see for example, [7], Chapter 2). Then,

$$\sum_{k>N} |h(k)| \leq \sum_{i=1}^n |r_i| \frac{|p_i|^N}{1 - |p_i|}. \tag{2}$$

Similar bounds are possible when the poles are not distinct.

The first purpose of this note is to present more sophisticated, and in many cases, substantially better bounds for the  $\ell^1$ -norm of the tail. These bounds are based on Theorem 2 of [2], which states that for the system (1),

$$|d| + \sigma_1(W_o^{\frac{1}{2}}W_c^{\frac{1}{2}}) \leq \|h\|_1 \leq |d| + 2 \sum_{i=1}^n \sigma_i(W_o^{\frac{1}{2}}W_c^{\frac{1}{2}}), \quad (3)$$

where

$$W_o = \sum_{k=0}^{\infty} (A^T)^k c^T c A^k \quad \text{and} \quad W_c = \sum_{k=0}^{\infty} A^k b b^T (A^T)^k$$

are the observability and controllability Gramians respectively [4].  $\sigma_i(W_o^{\frac{1}{2}}W_c^{\frac{1}{2}})$  are just the Hankel singular values of the system (1).

We now observe that  $\{0, h(N+1), h(N+2), \dots\}$ , the tail of the impulse response of system (1), is just the impulse response of the system  $\{A, A^N b, c, 0\}$ . Applying bounds (3) to this system, we have for any  $N \geq 0$ ,

$$\sigma_1(W_o^{\frac{1}{2}}A^N W_c^{\frac{1}{2}}) \leq \sum_{k>N} |h(k)| \leq 2 \sum_{i=1}^n \sigma_i(W_o^{\frac{1}{2}}A^N W_c^{\frac{1}{2}}). \quad (4)$$

Thus, we have upper and lower bounds for  $\|h\|_1$ :

$$S_N + \sigma_1(W_o^{\frac{1}{2}}A^N W_c^{\frac{1}{2}}) \leq \|h\|_1 \leq S_N + 2 \sum_{i=1}^n \sigma_i(W_o^{\frac{1}{2}}A^N W_c^{\frac{1}{2}}), \quad \forall N \geq 0. \quad (5)$$

The ratio between the upper and lower bounds for  $\|h\|_1$  in (4) is at most  $2n$ , whereas the ratio between the residue-expansion based upper bound (2) and any lower bound can be arbitrarily large.

We next show that with increasing  $N$ , the difference between the upper and lower bounds converges monotonically to zero.  $W_o$  satisfies the Lyapunov equation

$$A^T W_o A - W_o + c^T c = 0,$$

which implies that

$$(A^T)^k W_o A^k - (A^T)^{k-1} W_o A^{k-1} + (A^T)^{k-1} c^T c A^{k-1} = 0$$

for  $k = 1, 2, \dots$  Therefore,

$$(W_o^{\frac{1}{2}} A^k W_c^{\frac{1}{2}})^T (W_o^{\frac{1}{2}} A^k W_c^{\frac{1}{2}}) \leq (W_o^{\frac{1}{2}} A^{k-1} W_c^{\frac{1}{2}})^T (W_o^{\frac{1}{2}} A^{k-1} W_c^{\frac{1}{2}}), \quad k = 1, 2, \dots$$

This immediately means

$$\sigma_i(W_o^{\frac{1}{2}} A^k W_c^{\frac{1}{2}}) \leq \sigma_i(W_o^{\frac{1}{2}} A^{k-1} W_c^{\frac{1}{2}}), \quad i = 1, 2, \dots, n \text{ and } k = 1, 2, \dots,$$

from which it follows that the difference between the upper and lower bounds in (5) converges monotonically to zero with increasing  $N$ .

The above argument shows that all of the Hankel singular values of the impulse response of the ‘tail’ system  $\{A, A^N b, c, 0\}$  decrease monotonically (to zero, since the system is stable) as  $N \rightarrow \infty$ . In fact, we can say more: If we normalize the Hankel singular values by dividing them by the first one, the number of ‘normalized’ Hankel singular values that converge to nonzero values as  $N \rightarrow \infty$  equals the number of ‘dominant’ Jordan blocks of  $A$ , that is, the number of Jordan blocks of  $A$  which

- correspond to an eigenvalue of  $A$  with maximum magnitude, and
- which have the largest size among all Jordan blocks corresponding to an eigenvalue with maximum magnitude.

Thus, for large  $N$ , the number of significant terms in the sum  $\sum_{i=1}^n \sigma_i(W_o^{\frac{1}{2}} A^N W_c^{\frac{1}{2}})$  is just the ‘effective order’ of the tail system  $\{A, A^N b, c, 0\}$ .

Finally, we discuss informally a scheme for finding

$$N_{\min} = \min \left\{ N \left| 2 \sum_{i=1}^n \sigma_i(W_o^{\frac{1}{2}} A^N W_c^{\frac{1}{2}}) - \sigma_1(W_o^{\frac{1}{2}} A^N W_c^{\frac{1}{2}}) < \epsilon \right. \right\},$$

which is the smallest value of  $N$  for which the difference between the upper and lower bounds in (5) is less than  $\epsilon$ . As a preliminary step,  $W_c^{\frac{1}{2}}$  and  $W_o^{\frac{1}{2}}$  are computed. Then:

1. *We find the smallest positive integer  $M$  such that  $N_{\min} \leq 2^M$ .*

This is done iteratively where at the  $k$ th iteration, we form the matrix  $A^{2^k}$  by squaring  $A^{2^{k-1}}$  and check if

$$\sigma_1(W_o^{\frac{1}{2}} A^{2^k} W_c^{\frac{1}{2}}) + 2 \sum_{i=2}^n \sigma_i(W_o^{\frac{1}{2}} A^{2^k} W_c^{\frac{1}{2}}) < \epsilon,$$

and stop if the condition is satisfied. Clearly,  $M$  iterations are needed. Each iteration involves three  $n \times n$  matrix multiplies and one computation of singular values. For use in part (2), we store the matrices  $\{A, A^2, \dots, A^{2^M}\}$ .

2. *By a simple bisection,  $N_{\min}$  is then located in the set  $\{2^{M-1}, 2^{M-1} + 1, \dots, 2^M\}$ .*

We assume that  $M \geq 2$ , since computing  $N_{\min}$  is trivial otherwise. We start by forming  $\tilde{A} = A^{(2^{M-1} + 2^{M-2})}$  and checking if

$$\sigma_1(W_o^{\frac{1}{2}} \tilde{A} W_c^{\frac{1}{2}}) + 2 \sum_{i=2}^n \sigma_i(W_o^{\frac{1}{2}} \tilde{A} W_c^{\frac{1}{2}}) < \epsilon.$$

(Note that since  $A^{2^{M-1}}$  and  $A^{2^{M-2}}$  are both already available from step (1), and therefore this involves three  $n \times n$  matrix multiplies and one computation of singular values.) If the answer is yes, then  $N$  lies in the set  $\{2^{M-1}, 2^{M-1} + 1, \dots, 2^{M-1} + 2^{M-2}\}$ . Otherwise,  $N$  lies in the set  $\{2^{M-1} + 2^{M-2}, \dots, 2^M\}$ . By continuing this process (at most  $M - 1$  times) of halving the set where  $N$  lies, we may compute  $N_{\min}$  exactly.

Once  $N_{\min}$  is found,  $S_{N_{\min}}$  can be computed to give  $\|h\|_1$  to within an absolute accuracy of  $\epsilon$  (assuming infinite precision arithmetic; we have not considered the effects of data rounding here).

The exact determination of  $N_{\min}$  takes approximately  $6M$  matrix multiplies and  $2M$  computations of singular values. Forming  $S_{N_{\min}}$  takes about  $N_{\min}$  matrix-vector multiplies and  $N_{\min}$  vector-vector inner products. (Recall that  $2^{M-1} < N_{\min} \leq 2^M$ .) Since computing singular values is by far the most expensive of the above calculations, it might prove advantageous to not compute  $N_{\min}$  exactly, but to instead use an upper bound obtained by terminating the bisection in step (2) earlier. Computation may be further reduced by first balancing system (1), so that the Gramians  $W_c$  and  $W_o$  are diagonal and equal.

We note that for calculating the  $\mathbf{H}_{\infty}$ -norm of system (1) to within a relative accuracy  $\epsilon$ , there exist methods (see [1]) where the computational effort involved depends only on  $\epsilon$  and the state dimension  $n$ . However for determining  $\|h\|_1$  using the bounds in (5) to within an accuracy of  $\epsilon$

(relative or absolute), the number of computations depends on the system matrices  $A$ ,  $b$ ,  $c$  and  $d$  as well. We know of no way to overcome this deficiency.

### 3 Bounds for the worst-case $\ell^\infty$ -gain

We now combine the results of the previous section with results from [5] to derive bounds for the worst-case  $\ell^\infty$ -gain of discrete-time LTI systems with diagonal uncertainty. We consider the system shown in Figure 1:  $H$  is a stable discrete-time LTI plant.  $\Delta_1, \Delta_2, \dots, \Delta_m$  are scalar LTI perturbations that act on the system. Now, for some notation (indices  $i, j = 1, 2, \dots, m$ ):

- $\delta_i$  : Impulse response of perturbation  $\Delta_i$ .
- $h_{00}$  : Open-loop ( $\Delta = 0$ ) impulse response from  $w$  to  $z$ .
- $h_{i0}$  : Open-loop ( $\Delta = 0$ ) impulse response from  $w$  to  $y_i$ .
- $h_{0i}$  : Open-loop ( $\Delta = 0$ ) impulse response from  $u_i$  to  $z$ .
- $h_{cl}(\Delta)$  : Closed-loop impulse response from  $w$  to  $z$ .

We assume that  $\|\delta_i\|_1 \leq 1$  and denote by  $\Omega$  the corresponding set of all possible perturbations  $\Delta$ .

The quantity of interest is the worst-case (*i.e.* maximum possible)  $\ell^\infty$ -gain from  $w$  to  $z$ , which we define as

$$L_{wc} = \sup_{\Delta \in \Omega} \|h_{cl}(\Delta)\|_1.$$

In [5], Khammash and Pearson show that the  $L_{wc} \geq 1$  if and only if the following condition holds:

*There exists some nonzero  $x = [x_0, \dots, x_m]$  with  $x_i \geq 0$  such that*

$$x_i \leq \sum_{j=0}^m \|h_{ij}\|_1 x_j \quad i = 0, 1, \dots, m. \quad (\text{COND})$$

Condition (COND) may be expressed simply in terms of a matrix whose  $(i, j)$ -entry is  $\|h_{ij}\|_1$ ,  $i, j = 0, 1, \dots, m$ .

**Fact 1** *Condition COND holds if and only if the spectral radius of the matrix*

$$M = \begin{bmatrix} \|h_{00}\|_1 & \|h_{01}\|_1 & \cdots & \|h_{0m}\|_1 \\ \|h_{10}\|_1 & \|h_{11}\|_1 & \cdots & \|h_{1m}\|_1 \\ \vdots & \vdots & \ddots & \vdots \\ \|h_{m0}\|_1 & \|h_{m1}\|_1 & \cdots & \|h_{mm}\|_1 \end{bmatrix}$$

*is at least one.*

This fact, stated without proof in Theorem 1 of [6], is immediate from the following characterization of the spectral radius of a nonnegative matrix (a matrix with nonnegative entries)  $M$  (see, for example, page 504, corollary 8.3.3 of [3]):

$$\rho(M) = \max_{x \geq 0, x \neq 0} \min_{0 \leq i \leq m, x_i \neq 0} \frac{1}{x_i} \sum_{j=0}^m M_{ij} x_j.$$

( $M_{ij}$  refers to the  $(i, j)$ -entry of  $M$ .)

By simply scaling  $w$  and  $z$  by  $1/\sqrt{\gamma}$  ( $\gamma > 0$ ) as in Figure 2, and applying Fact 1, we conclude that

$$L_{\text{wc}} = \sup\{\gamma \mid \rho(D_\gamma M D_\gamma) \geq 1\}, \quad (6)$$

where

$$D_\gamma = \begin{bmatrix} 1/\sqrt{\gamma} & 0 \\ 0 & I \end{bmatrix}.$$

For convenience, we partition  $M$  as

$$M = \begin{bmatrix} M^{(11)} & M^{(12)} \\ M^{(21)} & M^{(22)} \end{bmatrix}, \quad (7)$$

where  $M^{(11)} \in \mathbf{R}_+$ ,  $M^{(12)} \in \mathbf{R}_+^{1 \times m}$ ,  $M^{(21)} \in \mathbf{R}_+^{m \times 1}$  and  $M^{(22)} \in \mathbf{R}_+^{m \times m}$ .

If  $\rho(D_\gamma M D_\gamma) \geq 1$  for all  $\gamma > 0$ , then we define  $L_{\text{wc}} = \infty$ . This corresponds to the case when  $\rho(M^{(22)}) \geq 1$ , and the system is not  $\ell^\infty$ -stable (see [5]). On the other hand, if  $\rho(D_\gamma M D_\gamma) < 1$  for all  $\gamma > 0$ , we define  $L_{\text{wc}} = 0$ . This corresponds to the case when either the first row (or the first column) of  $M$  is identically zero (with  $\rho(M^{(22)}) < 1$ ). Then  $h_{\text{cl}}(\Delta) = 0$  for all  $\Delta$ .

Of course, every entry of  $M$  is the  $\ell^\infty$ -gain of some LTI system; therefore, the remarks made in Section 2 about computing  $\ell^\infty$ -gains apply here as well. We may however use the fact that  $M$  is

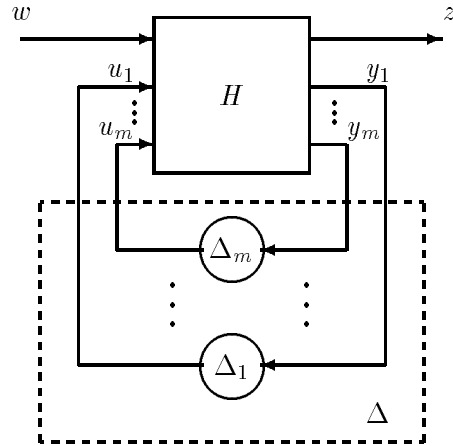


Figure 1: Linear system with diagonal uncertainty  $\Delta$ .

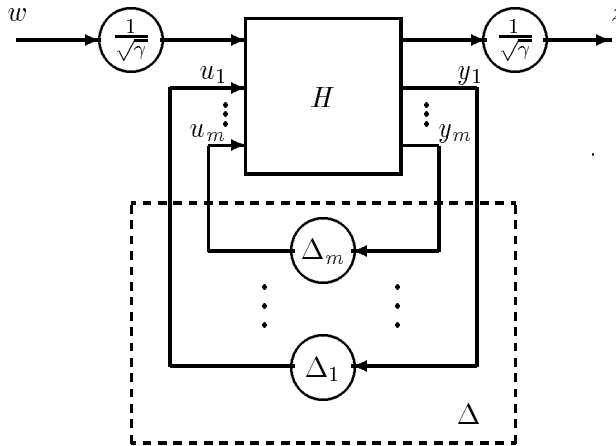


Figure 2: Uncertain linear system with the impulse response from  $w$  to  $z$  scaled by  $1/\gamma$ .



nonnegative to derive bounds on  $L_{\text{wc}}$  based on the bounds for the entries of  $M$ . We start with the following fact.

**Fact 2** *The spectral radius of a nonnegative matrix is a nondecreasing function of its entries.*

(See Corollary 8.1.19 on page 491 of [3].)

Fact 2 implies that  $\rho(D_\gamma P D_\gamma)$  is a nondecreasing function of the entries of the nonnegative matrix  $P$  and a nonincreasing function of  $\gamma > 0$ . These, in turn, mean that the function  $\Phi(P)$  of a nonnegative matrix  $P$  defined by

$$\Phi(P) = \sup\{\gamma \mid \rho(D_\gamma P D_\gamma) \geq 1\}$$

is nondecreasing with the entries of  $P$ . We then have the following bounds for  $L_{\text{wc}}$ :

**Theorem 1** *Let  $\alpha_{ij}^N$  and  $\beta_{ij}^N$  be lower and upper bounds for  $\|h_{ij}\|_1$  computed via equation (5) for some  $N > 0$ . Let  $M_{\text{lb}}^N$  and  $M_{\text{ub}}^N$  be matrices with  $(i, j)$ -entry  $\alpha_{ij}^N$  and  $\beta_{ij}^N$  respectively ( $i, j = 0, 1, \dots, m$ ). Then*

$$L_{\text{lb}}^N = \Phi(M_{\text{lb}}^N) = \sup\{\gamma \mid \rho(D_\gamma M_{\text{lb}}^N D_\gamma) \geq 1\},$$

and

$$L_{\text{ub}}^N = \Phi(M_{\text{ub}}^N) = \sup\{\gamma \mid \rho(D_\gamma M_{\text{ub}}^N D_\gamma) \geq 1\},$$

are lower and upper bounds respectively for  $L_{\text{wc}}$ , i.e.  $L_{\text{lb}}^N \leq L_{\text{wc}} \leq L_{\text{ub}}^N$ .

Computation of  $L_{\text{lb}}^N$  and  $L_{\text{ub}}^N$  is straightforward, once we make the following observation:

**Fact 3** *The spectral radius of a nonnegative matrix is also an eigenvalue.*

(See Theorem 8.3.1 on page 503 of [3].)

Given a  $(m+1) \times (m+1)$  matrix  $P$ , we first partition conformally as with  $M$  in equation (7) as

$$P = \begin{bmatrix} P^{(11)} & P^{(12)} \\ P^{(21)} & P^{(22)} \end{bmatrix}.$$

Then,  $\Phi(P) = \infty$  if  $\rho(P^{(22)}) \geq 1$ . Otherwise, we note that  $\rho(D_\gamma P D_\gamma) = \rho(D_\gamma^2 P)$ , and solve for  $D_\gamma^2 P x = x$  for some nonzero  $(m+1)$ -vector  $x$  to obtain

$$\Phi(P) = P^{(11)} + P^{(12)}(I - P^{(22)})^{-1}P^{(21)}.$$

The above formula shows that if  $\rho(P^{(22)}) < 1$ ,  $\Phi(P)$  is just the unique solution to the equation  $\rho(D_\gamma P D_\gamma) = 1$ .

## 4 Conclusion

We have presented simple bounds on the  $\ell^\infty$ -gain of single-input single-output linear discrete time systems. We have shown how to combine these bounds with recent results from [5] to compute guaranteed bounds for the worst-case  $\ell^\infty$  gain of discrete-time LTI systems with diagonal uncertainty. The bounds may be easily extended to block diagonal uncertainties as well as to continuous time systems.

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