

Mathematical Logic

Part One

Announcements

- Problem Session tonight from 7:00 – 7:50 in 380-380X.
 - Optional, but highly recommended!
- Problem Set 3 Checkpoint due right now.
- 2× Handouts
 - Problem Set 3 Checkpoint Solutions
 - Diagonalization
- Problem Set 2 Solutions distributed at end of class.

Office Hours

- We finally have stable office hours locations!
- Website will be updated soon with details.

An Important Question

How do we formalize the logic we've been using in our proofs?

Where We're Going

- **Propositional Logic** (Today)
 - Basic logical connectives.
 - Truth tables.
 - Logical equivalences.
- **First-Order Logic** (Today / Wednesday)
 - Reasoning about properties of multiple objects.

Propositional Logic

A **proposition** is a statement that is,
by itself, either true or false.

Some Sample Propositions

- Puppies are cuter than kittens.
- Kittens are cuter than puppies.
- Usain Bolt can outrun everyone in this room.
- CS103 is useful for cocktail parties.
- This is the last entry on this list.

More Propositions

- I'm a single lady.
- This place about to blow.
- Party rock is in the house tonight.
- We can dance if we want to.
- We can leave your friends behind.

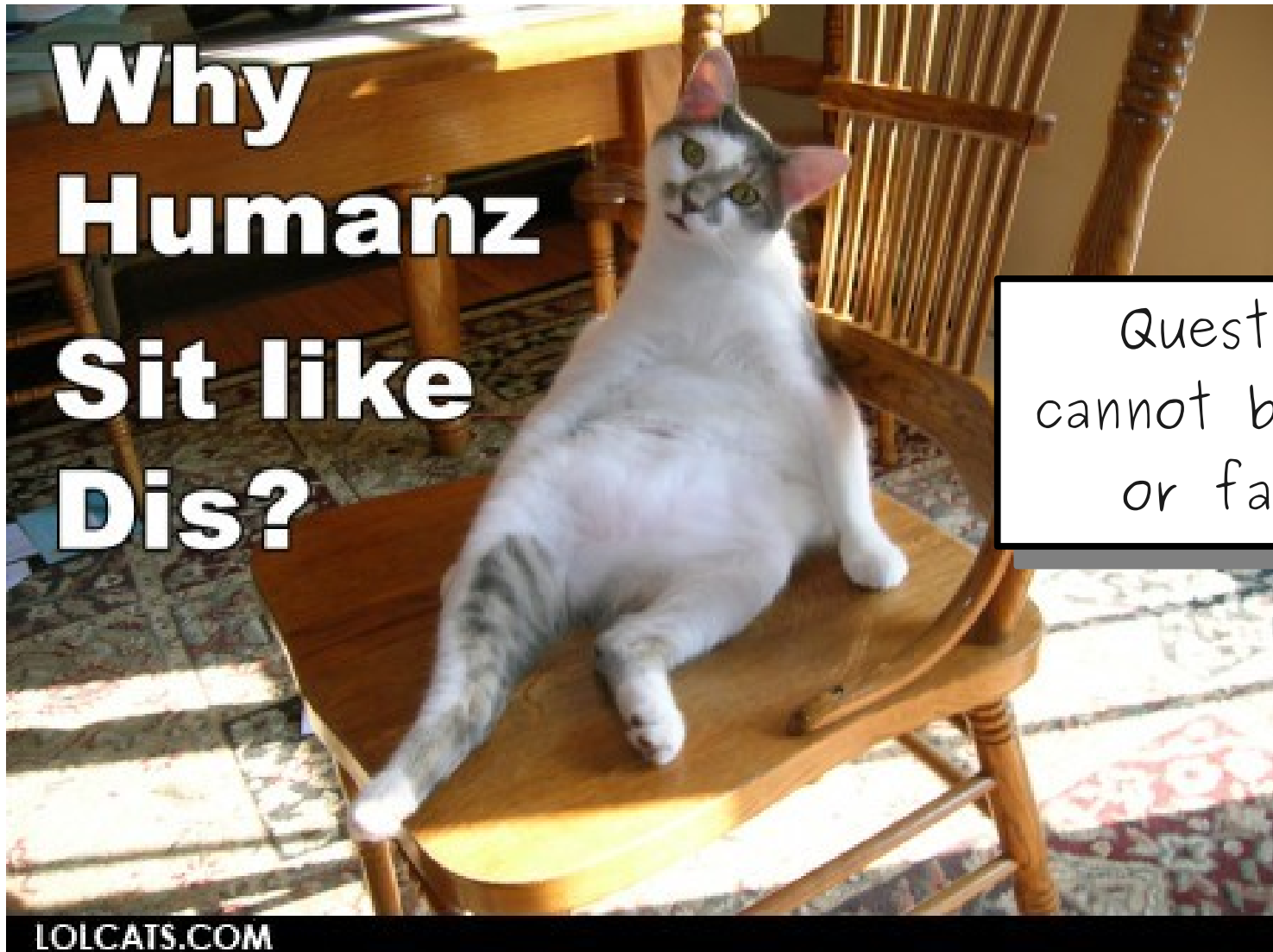
Things That Aren't Propositions



Commands
cannot be true
or false.

FLY, YOU FOOLS!

Things That Aren't Propositions

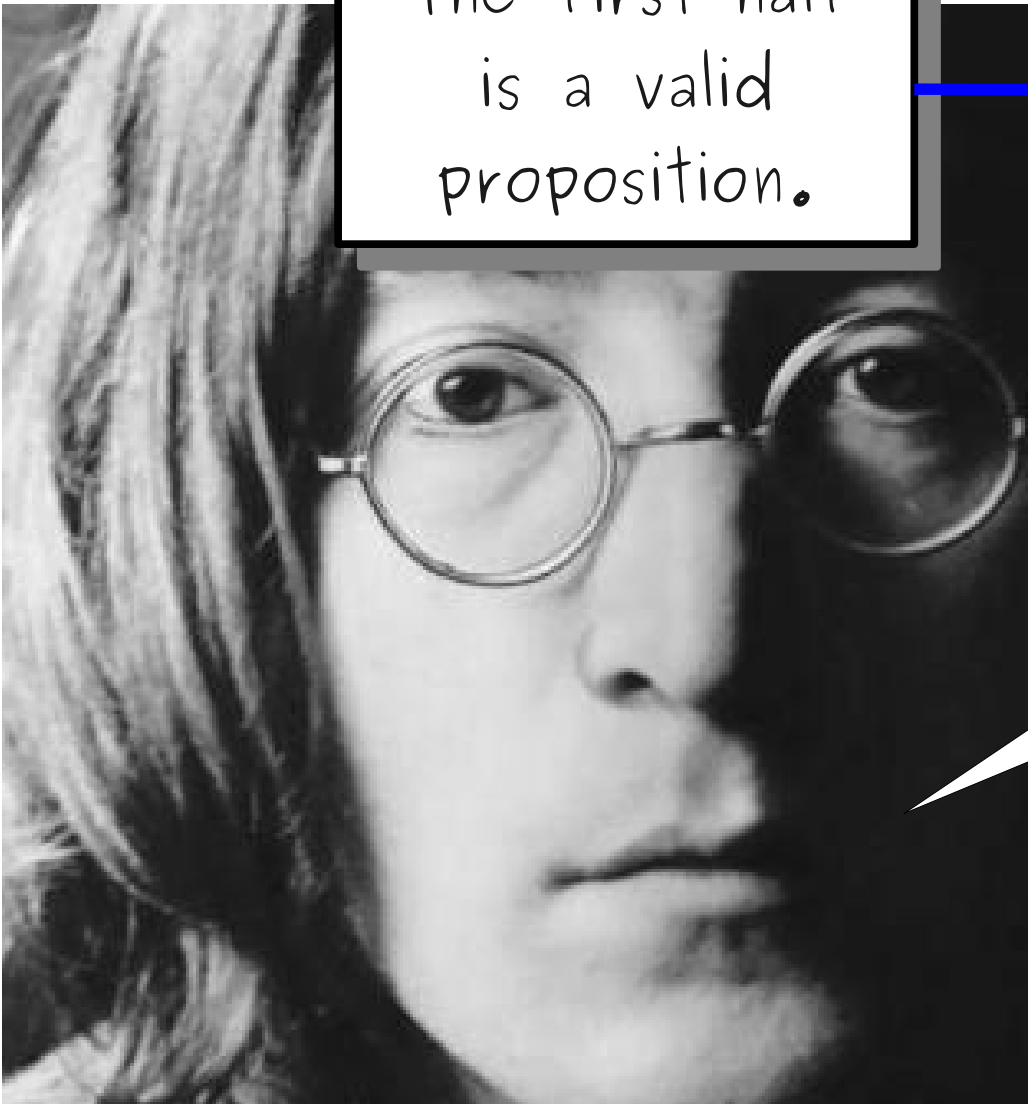


**Why
Humanz
Sit like
Dis?**

Questions
cannot be true
or false.

LOLCATS.COM

Things That Aren't Propositions



The first half
is a valid
proposition.

I am the walrus,
goo goo g'joob

Jibberish cannot
be true or
false.

Propositional Logic

- **Propositional logic** is a mathematical system for reasoning about propositions and how they relate to one another.
- Propositional logic enables us to
 - Formally encode how the truth of various propositions influences the truth of other propositions.
 - Determine if certain combinations of propositions are always, sometimes, or never true.
 - Determine whether certain combinations of propositions logically entail other combinations.

Variables and Connectives

- Propositional logic is a formal mathematical system whose syntax is rigidly specified.
- Every statement in propositional logic consists of **propositional variables** combined via **logical connectives**.
 - Each variable represents some proposition, such as “You wanted it” or “You should have put a ring on it.”
 - Connectives encode how propositions are related, such as “If you wanted it, you should have put a ring on it.”

Propositional Variables

- Each proposition will be represented by a **propositional variable**.
- Propositional variables are usually represented as lower-case letters, such as p , q , r , s , etc.
 - If we need more, we can use subscripts: p_1 , p_2 , etc.
- Each variable can take one of two values: true or false.

Logical Connectives

- **Logical NOT: $\neg p$**
 - Read “**not** p ”
 - $\neg p$ is true if and only if p is false.
 - Also called **logical negation**.
- **Logical AND: $p \wedge q$**
 - Read “ p **and** q .”
 - $p \wedge q$ is true if both p and q are true.
 - Also called **logical conjunction**.
- **Logical OR: $p \vee q$**
 - Read “ p **or** q .”
 - $p \vee q$ is true if at least one of p or q are true (inclusive OR)
 - Also called **logical disjunction**.

Truth Tables

p	q	$p \wedge q$
F	F	F
F	T	F
T	F	F
T	T	T

Truth Tables

p	q	$p \wedge q$
F	F	F
F	T	F
T	F	F
T	T	T

If p is false and q is false, then "both p and q " is false.

Truth Tables

p	q	$p \wedge q$
F	F	F
F	T	F
T	F	F
T	T	T

Truth Tables

p	q	$p \wedge q$
F	F	F
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Truth Tables

p	q	$p \wedge q$
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Truth Tables

p	q	$p \wedge q$
F	F	F
F	T	F
T	F	F
T	T	T

"both p and q " is true only when both p and q are true.



Truth Tables

Truth Tables

p	q	$p \vee q$
F	F	F
F	T	T
T	F	T
T	T	T

Truth Tables

p	q	$p \vee q$
F	F	F
F	T	T
T	F	T
T	T	T

"p or q" is true even if both p and q are true.

Remember that there are three ways for "p or q" to be true!

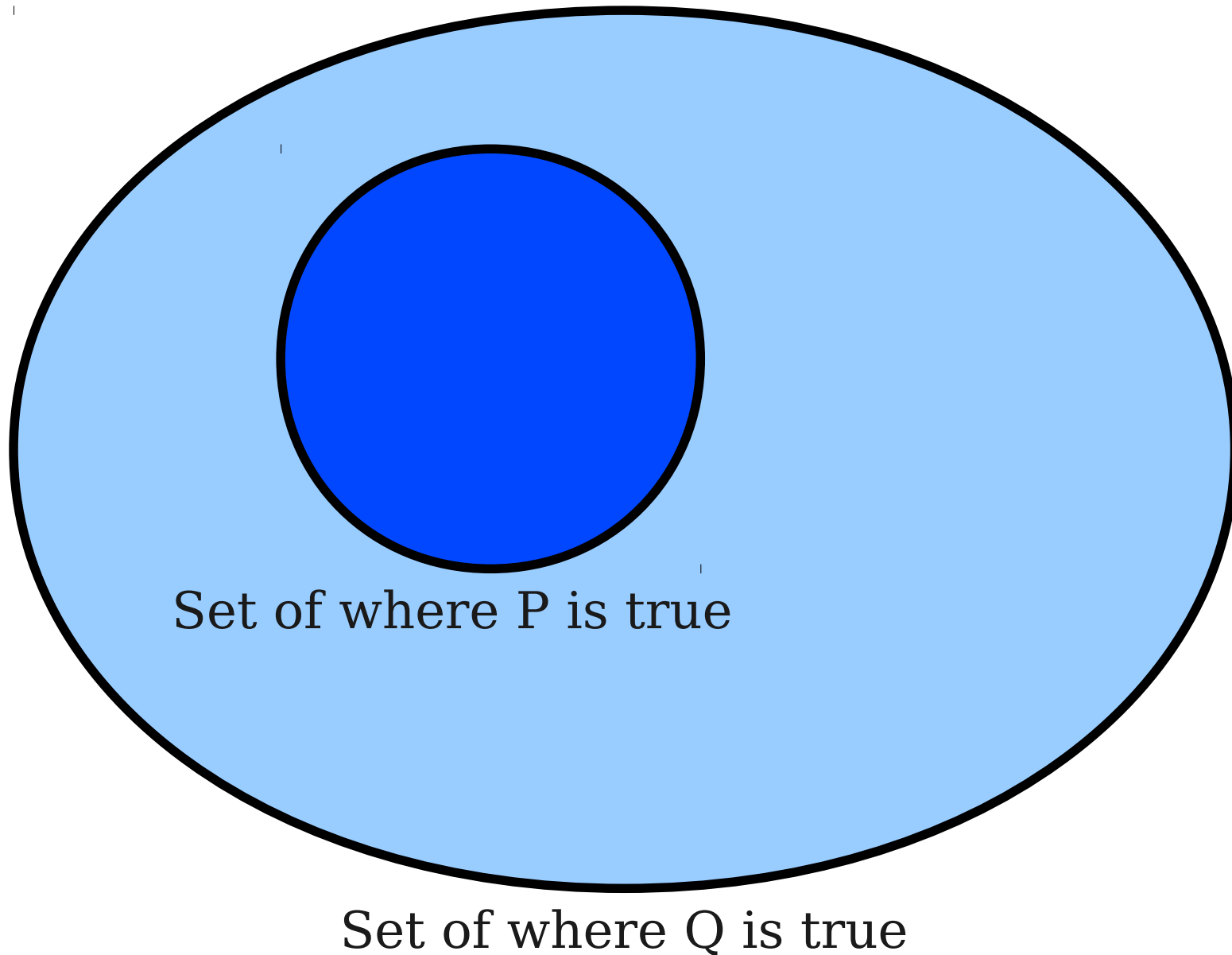
Truth Tables

p	$\neg p$
F	T
T	F

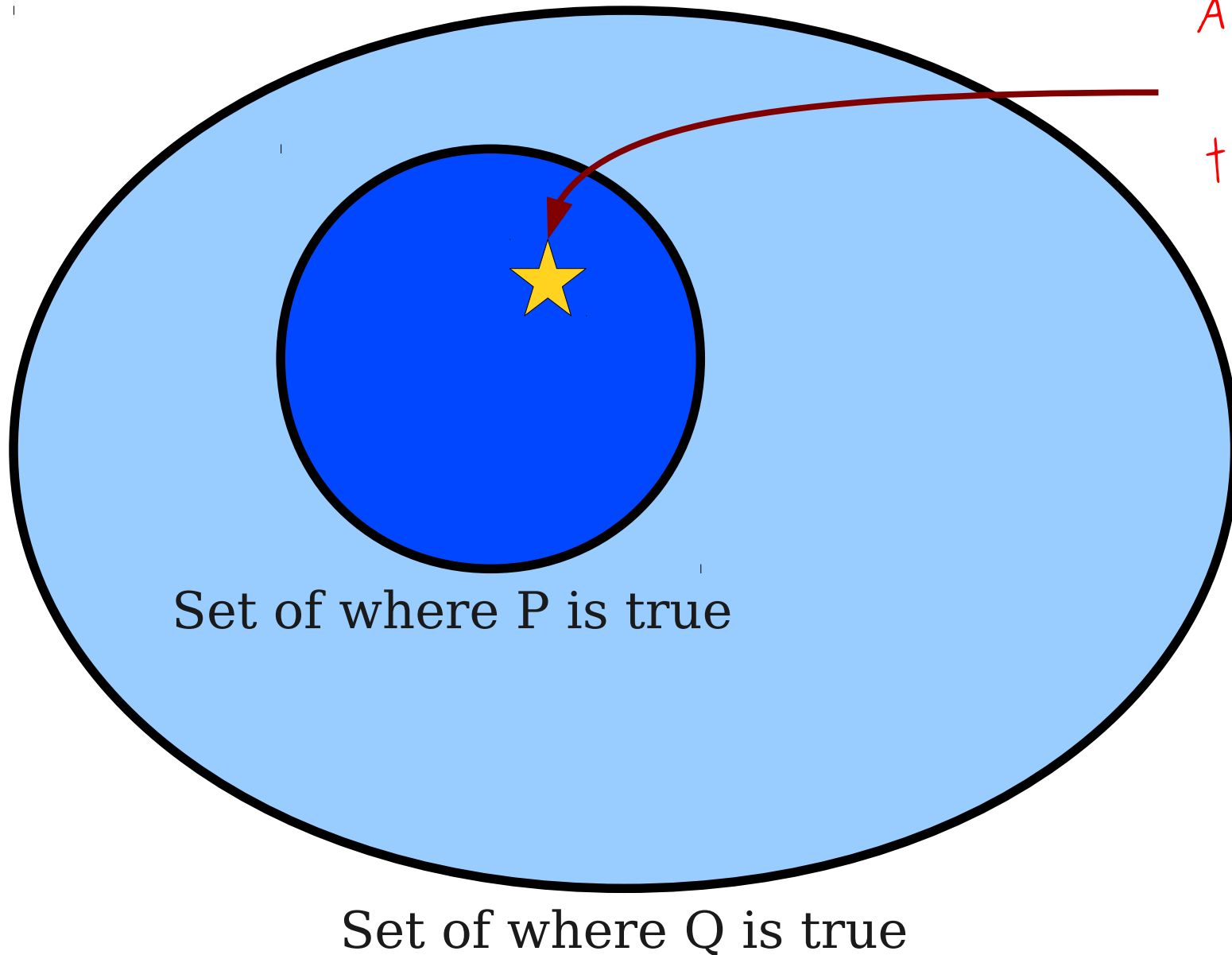
Implication

- An important connective is logical implication: $p \rightarrow q$.
- Recall: $p \rightarrow q$ means “if p is true, q is true as well.”
- Recall: $p \rightarrow q$ says **nothing** about what happens if p is false.
- Recall: $p \rightarrow q$ says **nothing** about causality; it just says that if p is true, q will be true as well.

Implication, Diagrammatically

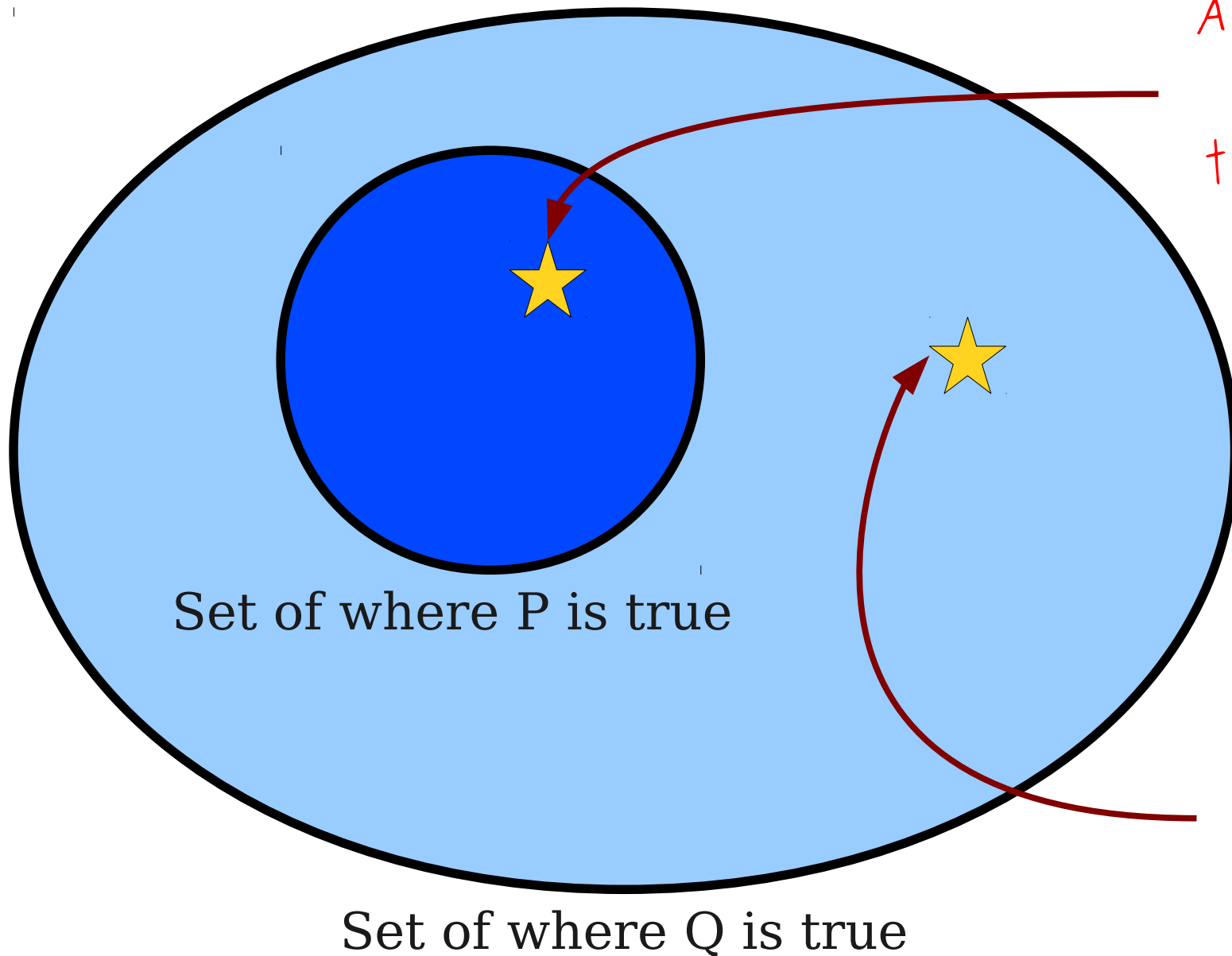


Implication, Diagrammatically



Any time P is true, Q is true as well.

Implication, Diagrammatically



Any time P is true, Q is true as well.

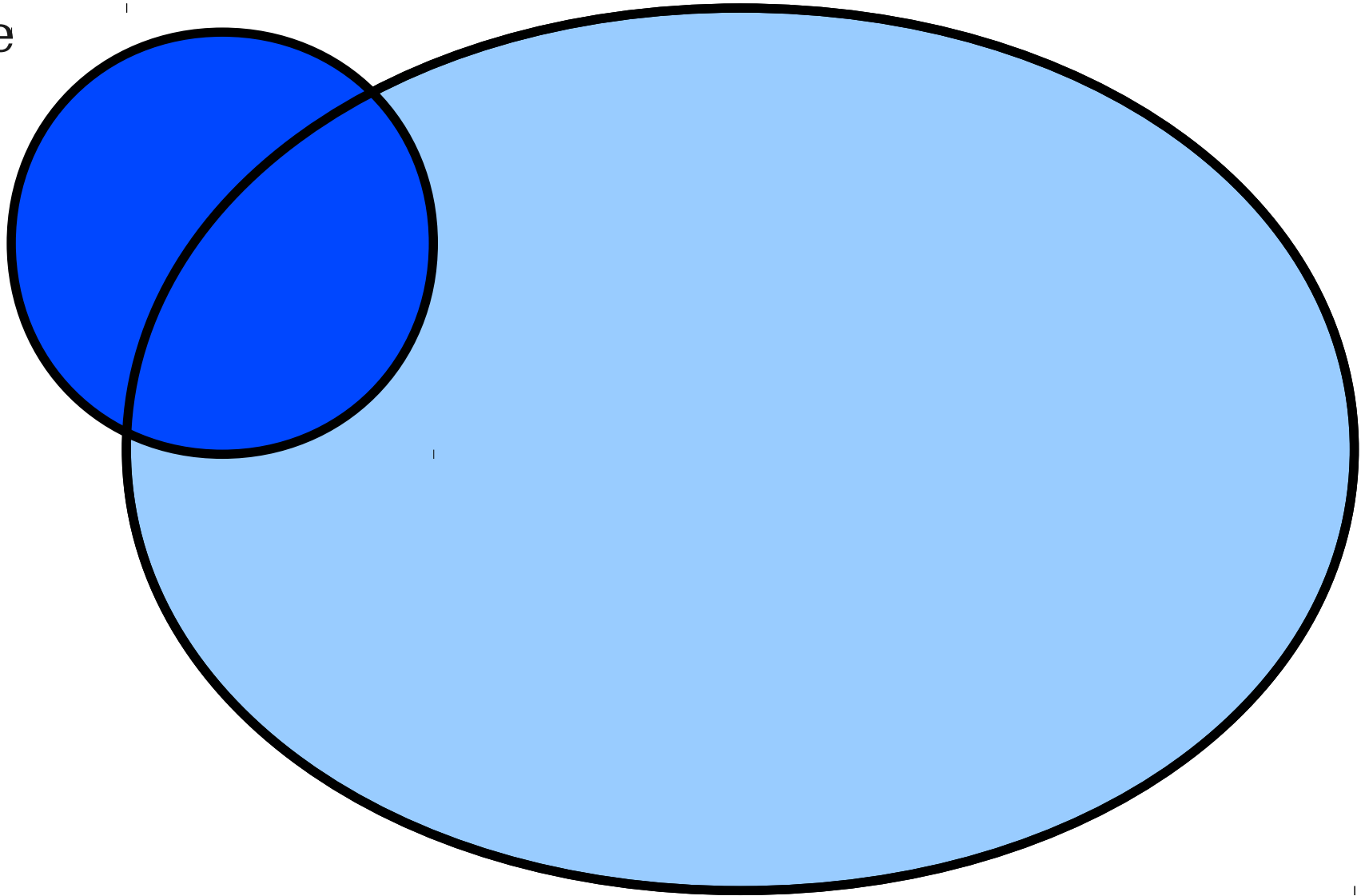
Any time P isn't true, Q may or may not be true.

When p Does Not Imply q

- $p \rightarrow q$ means “if p is true, q is true as well.”
- Recall: The **only way** for $p \rightarrow q$ to be false is if we know that p is true but q is false.
- Rationale:
 - If p is false, $p \rightarrow q$ doesn't guarantee anything. It's true, but it's not **meaningful**.
 - If p is true and q is true, then the statement “if p is true, then q is also true” is itself true.
 - If p is true and q is false, then the statement “if p is true, q is also true” is false.

$P \rightarrow Q$ is false

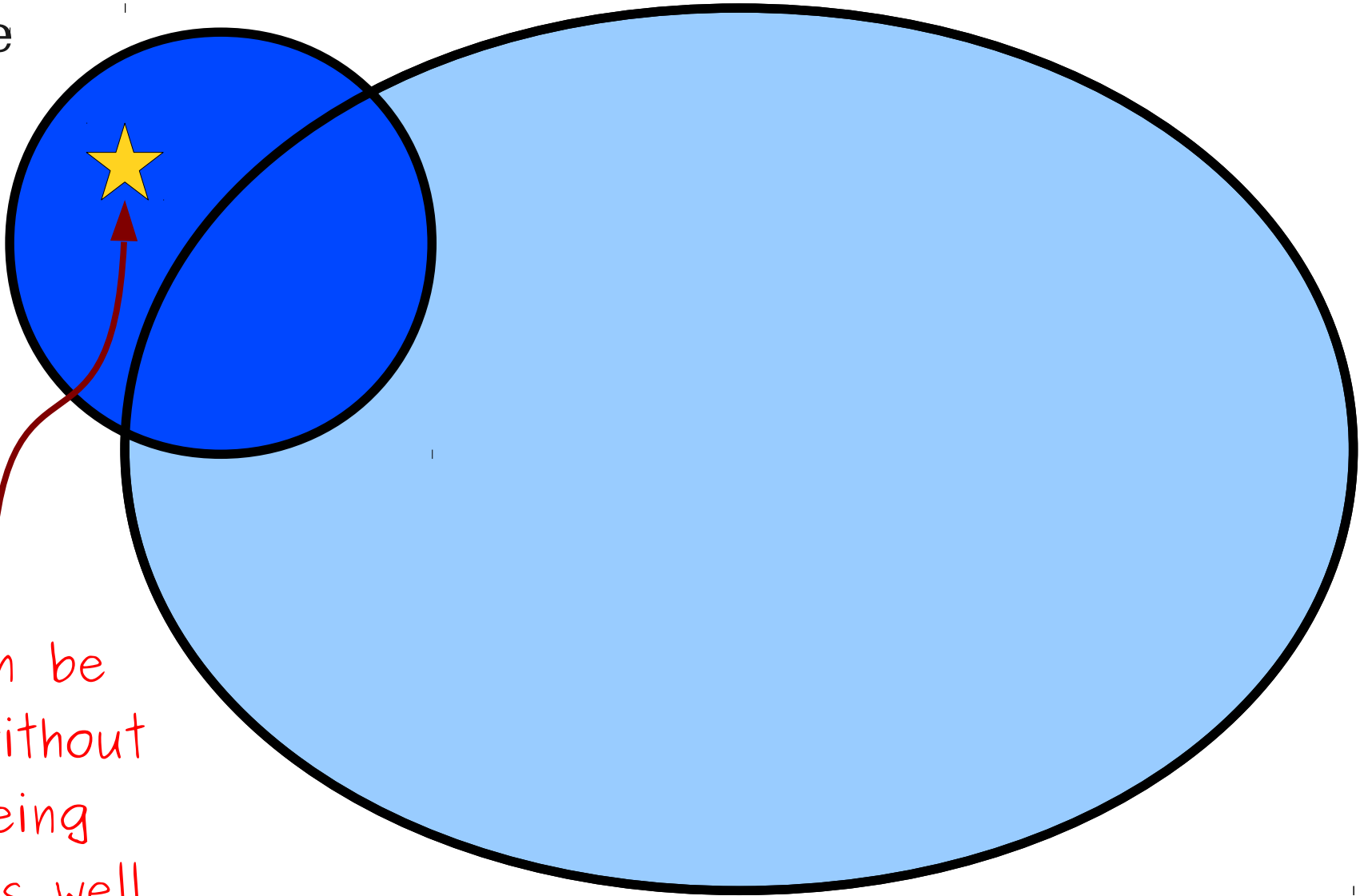
Set of
where
P is
true



Set of where Q is true

$P \rightarrow Q$ is false

Set of
where
P is
true



P can be
true without
Q being
true as well

Set of where Q is true

Truth Table for Implication

p	q	$p \rightarrow q$
F	F	
F	T	
T	F	
T	T	

Truth Table for Implication

p	q	$p \rightarrow q$
F	F	
F	T	
T	F	
T	T	

Truth Table for Implication

p	q	$p \rightarrow q$
F	F	
F	T	
T	F	
T	T	

In both of these cases,
 p is false, so the
statement "if p , then q "
is vacuously true.

Truth Table for Implication

p	q	$p \rightarrow q$
F	F	T
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T	F	
T	T	

Truth Table for Implication

p	q	$p \rightarrow q$
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F	T	T
T	F	
T	T	

$p \rightarrow q$ should mean
when p is true, q is
true as well. But here
 p is true and q is
false!

Truth Table for Implication

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Truth Table for Implication

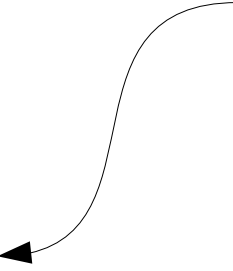
p	q	$p \rightarrow q$
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$p \rightarrow q$ means that if we ever find that p is true, we'll find that q is true as well.

Truth Table for Implication

p	q	$p \rightarrow q$
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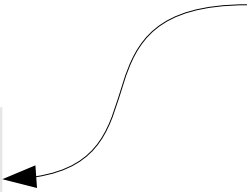
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Truth Table for Implication

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The only way for $p \rightarrow q$ to be false is for p to be true and q to be false.



The Biconditional

- The **biconditional** connective $p \leftrightarrow q$ is read “ p if and only if q .”
- Intuitively, either both p and q are true, or neither of them are.

p	q	$p \leftrightarrow q$
F	F	
F	T	
T	F	
T	T	

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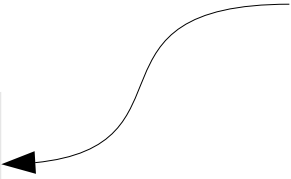
One of p or q is true without the other.

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Both p and q are false here, so the statement “ p if and only if q ” is true.

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One interpretation of \leftrightarrow is to think of it as equality: the two propositions must have equal truth values.

True and False

- There are two more “connectives” to speak of: true and false.
- The symbol \top is a value that is always true.
- The symbol \perp is value that is always false.
- These are often called connectives, though they don't connect anything.
 - (Or rather, they connect zero things.)

Operator Precedence

- How do we parse this statement?

$$\neg x \rightarrow y \vee z \rightarrow x \vee y \wedge z$$

- Operator precedence for propositional logic:

\neg

\wedge

\vee

\rightarrow

\leftrightarrow

- All operators are right-associative.
- We can use parentheses to disambiguate.

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Operator Precedence

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$$(\neg x) \rightarrow y \vee z \rightarrow x \vee y \wedge z$$

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\rightarrow

\leftrightarrow

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Operator Precedence

- How do we parse this statement?

$$(\neg x) \rightarrow y \vee z \rightarrow x \vee (y \wedge z)$$

- Operator precedence for propositional logic:

\neg

\wedge

\vee

\rightarrow

\leftrightarrow

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Operator Precedence

- How do we parse this statement?

$$(\neg x) \rightarrow y \vee z \rightarrow x \vee (y \wedge z)$$

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\neg

\wedge

\vee

\rightarrow

\leftrightarrow

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Operator Precedence

- How do we parse this statement?

$$(\neg x) \rightarrow (y \vee z) \rightarrow (x \vee (y \wedge z))$$

- Operator precedence for propositional logic:

\neg

\wedge

\vee

\rightarrow

\leftrightarrow

- All operators are right-associative.
- We can use parentheses to disambiguate.

Operator Precedence

- How do we parse this statement?

$$(\neg x) \rightarrow (y \vee z) \rightarrow (x \vee (y \wedge z))$$

- Operator precedence for propositional logic:

\neg

\wedge

\vee

\rightarrow

\leftrightarrow

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$$(\neg x) \rightarrow ((y \vee z) \rightarrow (x \vee (y \wedge z)))$$

- Operator precedence for propositional logic:

\neg

\wedge

\vee

\rightarrow

\leftrightarrow

- All operators are right-associative.
- We can use parentheses to disambiguate.

Recap So Far

- A **propositional variable** is a variable that is either true or false.
- The **logical connectives** are
 - Negation: $\neg p$
 - Conjunction: $p \wedge q$
 - Disjunction: $p \vee q$
 - Implication: $p \rightarrow q$
 - Biconditional: $p \leftrightarrow q$
 - True: \top
 - False: \perp

Translating into Propositional Logic

Some Sample Propositions

a: There is a velociraptor outside my apartment.

b: Velociraptors can open windows.

c: I am in my apartment right now.

d: My apartment has windows.

e: I am going to be eaten by a velociraptor

I won't be eaten by a velociraptor if there isn't a velociraptor outside my apartment.

$$\neg a \rightarrow \neg e$$

“ p if q ”

translates to

$$q \rightarrow p$$

It does **not** translate to

$$p \rightarrow q$$

Some Sample Propositions

a: There is a velociraptor outside my apartment.

b: Velociraptors can open windows.

c: I am in my apartment right now.

d: My apartment has windows.

e: I am going to be eaten by a velociraptor

If there is a velociraptor outside my apartment, but it can't open windows, I am not going to be eaten by a velociraptor.

$$a \wedge \neg b \rightarrow \neg e$$

“ p , but q ”

translates to

$p \wedge q$

Some Sample Propositions

a: There is a velociraptor outside my apartment.

b: Velociraptors can open windows.

c: I am in my apartment right now.

d: My apartment has windows.

e: I am going to be eaten by a velociraptor

I am only in my apartment when there are no velociraptors outside.

$$c \rightarrow \neg a$$

“ p only when q ”

translates to

$$p \rightarrow q$$

The Takeaway Point

- When translating into or out of propositional logic, be very careful not to get tripped up by nuances of the English language.
 - In fact, this is one of the reasons we have a symbolic notation in the first place!
- Many prepositions lead to counterintuitive translations; make sure to double-check yourself!

Logical Equivalence

More Elaborate Truth Tables

p	q	$p \wedge (p \rightarrow q)$
F	F	
F	T	
T	F	
T	T	

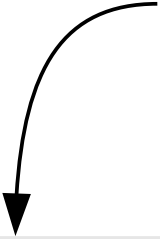
More Elaborate Truth Tables

We can't evaluate this until we have a value for $p \rightarrow q$.

p	q	$p \wedge (p \rightarrow q)$
F	F	
F	T	
T	F	
T	T	

More Elaborate Truth Tables

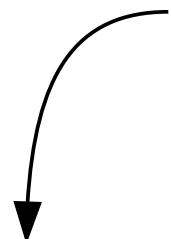
so let's start by evaluating this.



p	q	$p \wedge (p \rightarrow q)$
F	F	
F	T	
T	F	
T	T	

More Elaborate Truth Tables

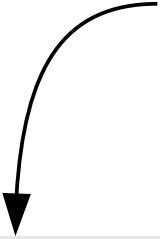
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p	q	$p \wedge (p \rightarrow q)$
F	F	
F	T	
T	F	
T	T	

More Elaborate Truth Tables

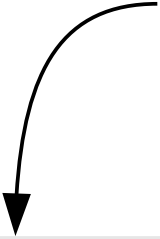
so let's start by evaluating this.



p	q	$p \wedge (p \rightarrow q)$
F	F	T
F	T	
T	F	
T	T	

More Elaborate Truth Tables

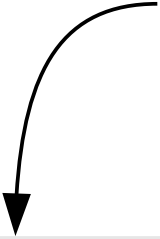
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p	q	$p \wedge (p \rightarrow q)$
F	F	T
F	T	
T	F	
T	T	

More Elaborate Truth Tables

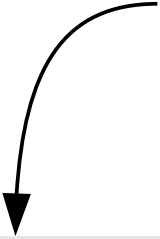
so let's start by evaluating this.



p	q	$p \wedge (p \rightarrow q)$
F	F	T
F	T	T
T	F	
T	T	

More Elaborate Truth Tables

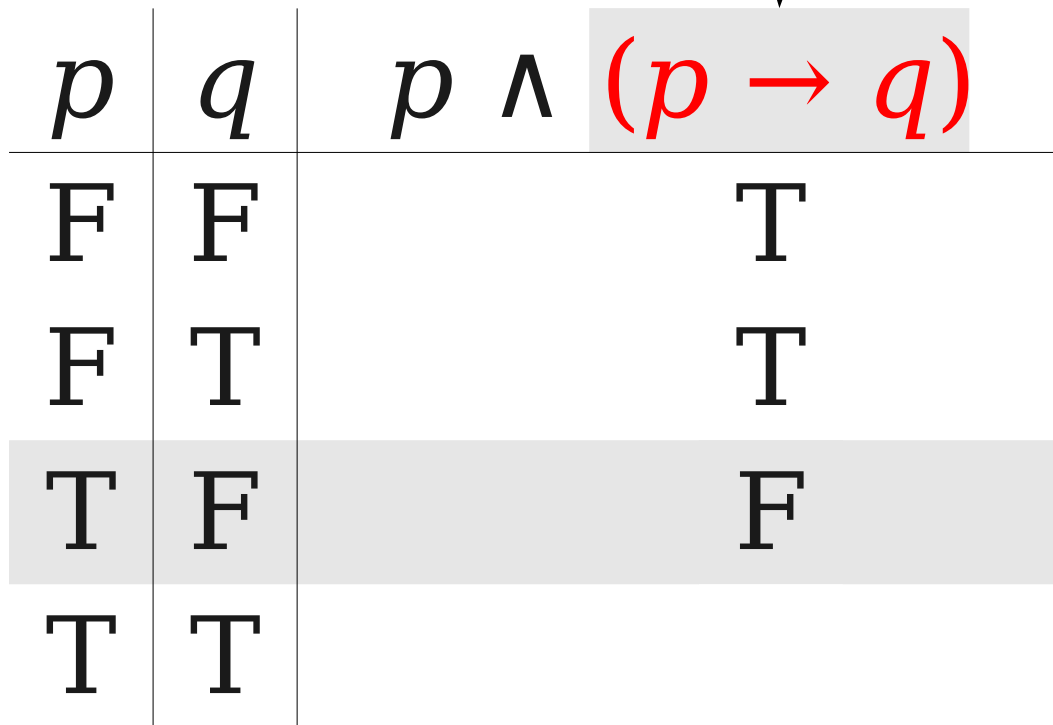
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p	q	$p \wedge (p \rightarrow q)$
F	F	T
F	T	T
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More Elaborate Truth Tables

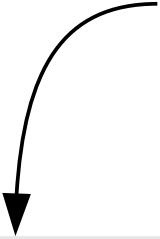
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More Elaborate Truth Tables

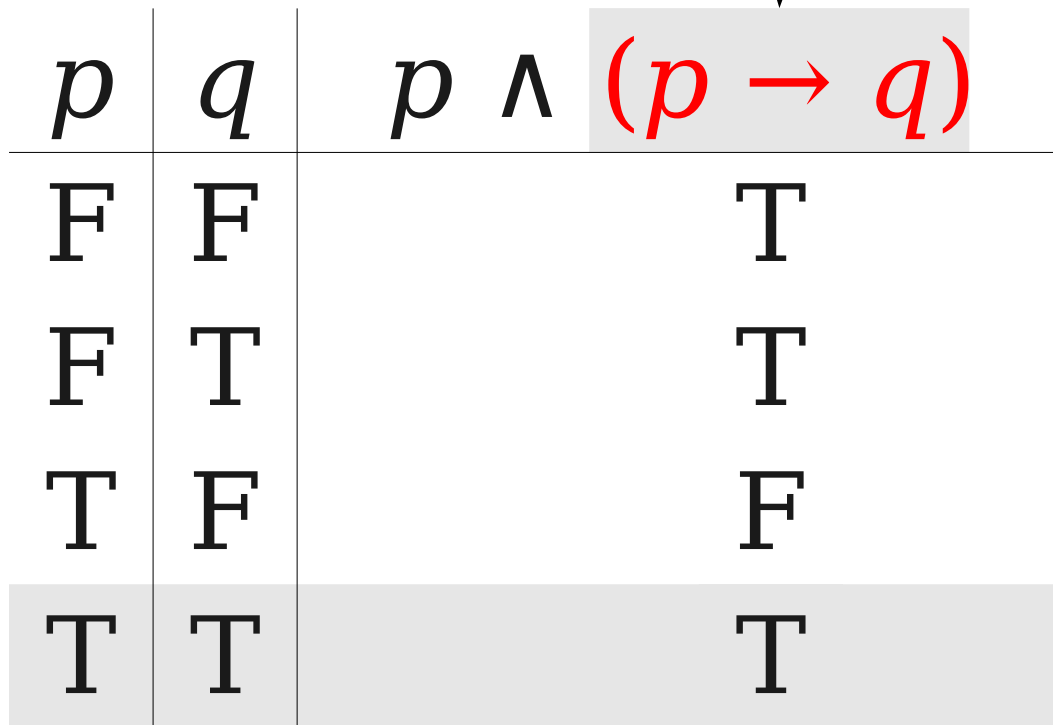
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p	q	$p \wedge (p \rightarrow q)$
F	F	T
F	T	T
T	F	F
T	T	

More Elaborate Truth Tables

so let's start by evaluating this.



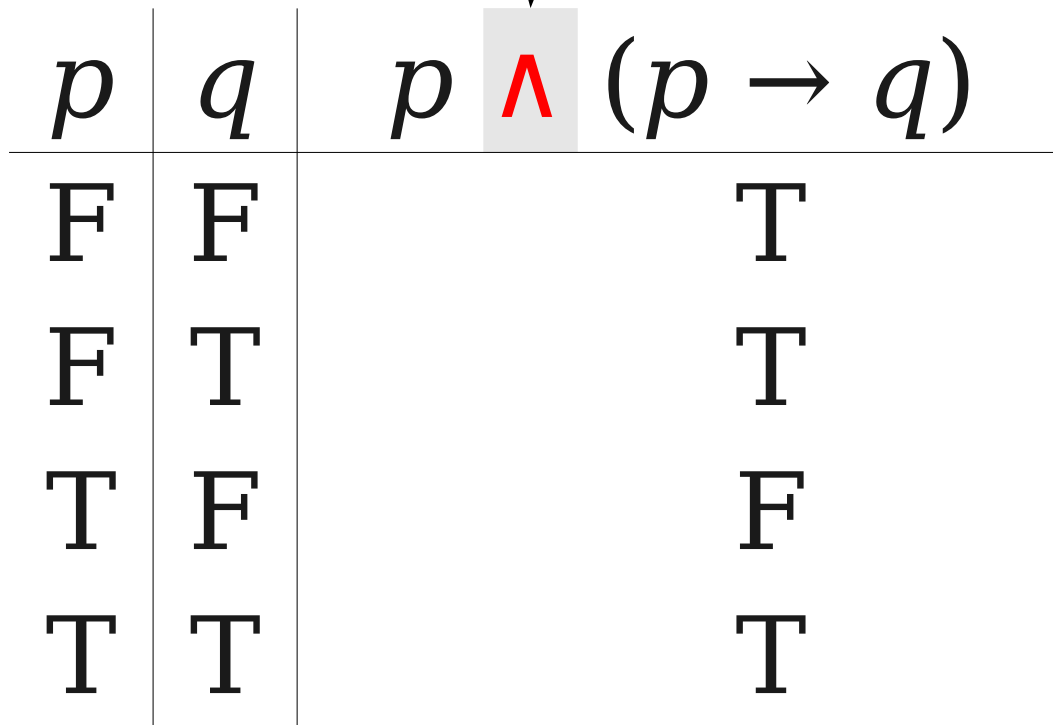
p	q	$p \wedge (p \rightarrow q)$
F	F	T
F	T	T
T	F	F
T	T	T

More Elaborate Truth Tables

p	q	$p \wedge (p \rightarrow q)$
F	F	T
F	T	T
T	F	F
T	T	T

More Elaborate Truth Tables

Now we can go evaluate this.



p	q	p	\wedge	$(p \rightarrow q)$
F	F			T
F	T			T
T	F			F
T	T			T

More Elaborate Truth Tables

Now we can go evaluate this.

p	q	$p \wedge (p \rightarrow q)$
F	F	T
F	T	T
T	F	F
T	T	T

More Elaborate Truth Tables

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p	q	$p \wedge (p \rightarrow q)$
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More Elaborate Truth Tables

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p	q	p	\wedge	$(p \rightarrow q)$
F	F	F		T
F	T	F		T
T	F			F
T	T			T

More Elaborate Truth Tables

Now we can go evaluate this.

p	q	p	\wedge	$(p \rightarrow q)$
F	F	F	F	T
F	T	F	F	T
T	F	T	F	F
T	T	T	T	T

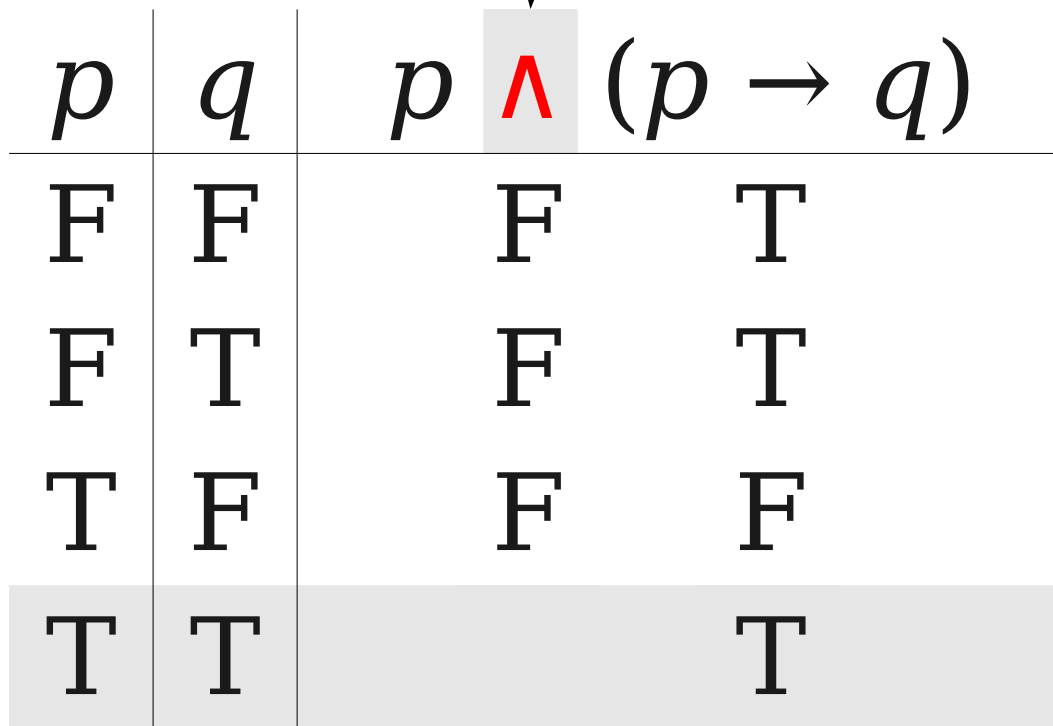
More Elaborate Truth Tables

Now we can go evaluate this.

p	q	p	\wedge	$(p \rightarrow q)$
F	F	F	F	T
F	T	F	F	T
T	F	F	F	F
T	T			T

More Elaborate Truth Tables

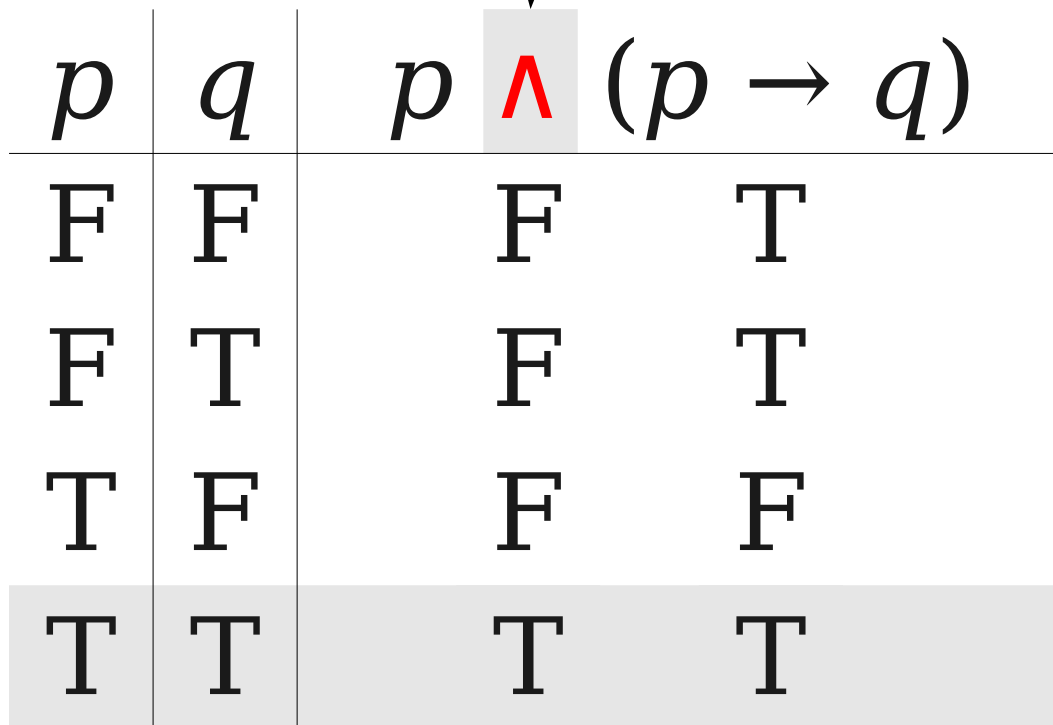
Now we can go evaluate this.



p	q	p	\wedge	$(p \rightarrow q)$
F	F	F	F	T
F	T	F	F	T
T	F	F	F	F
T	T	T	T	T

More Elaborate Truth Tables

Now we can go evaluate this.



p	q	p	\wedge	$(p \rightarrow q)$
F	F	F	F	T
F	T	F	F	T
T	F	F	F	F
T	T	T	T	T

More Elaborate Truth Tables

p	q	$p \wedge (p \rightarrow q)$
F	F	F T
F	T	F T
T	F	F F
T	T	T T

More Elaborate Truth Tables

This gives the final truth value for the expression.

p	q	$p \wedge (p \rightarrow q)$
F	F	F
F	T	F
T	F	F
T	T	T

Negations

- $p \wedge q$ is false if and only if $\neg(p \wedge q)$ is true.
- Intuitively, this is only possible if either p is false or q is false (or both!)
- In propositional logic, we can write this as $\neg p \vee \neg q$.
- How would we prove that $\neg(p \wedge q)$ and $\neg p \vee \neg q$ are equivalent?
- **Idea:** Build truth tables for both expressions and confirm that they always agree.

Negating AND

p	q	$\neg(p \wedge q)$
F	F	
F	T	
T	F	
T	T	

Negating AND

p	q	$\neg(p \wedge q)$
F	F	F
F	T	F
T	F	F
T	T	T

Negating AND

p	q	$\neg(p \wedge q)$
F	F	T
F	T	T
T	F	T
T	T	F

Negating AND

p	q	$\neg(p \wedge q)$
F	F	T
F	T	T
T	F	T
T	T	F

Negating AND

p	q	$\neg(p \wedge q)$
F	F	T
F	T	T
T	F	T
T	T	F

p	q	$\neg p \vee \neg q$
F	F	T
F	T	T
T	F	T
T	T	F

Negating AND

p	q	$\neg(p \wedge q)$
F	F	T
F	T	T
T	F	T
T	T	F

p	q	$\neg p \vee \neg q$
F	F	T
F	T	T
T	F	F
T	T	F

Negating AND

p	q	$\neg(p \wedge q)$
F	F	T
F	T	T
T	F	T
T	T	F

p	q	$\neg p \vee \neg q$
F	F	T
F	T	T
T	F	T
T	T	F

Negating AND

p	q	$\neg(p \wedge q)$
F	F	T
F	T	T
T	F	T
T	T	F

p	q	$\neg p$	\vee	$\neg q$
F	F	T	T	T
F	T	T	T	F
T	F	F	T	T
T	T	F	F	F

Negating AND

p	q	$\neg(p \wedge q)$
F	F	T
F	T	T
T	F	T
T	T	F

p	q	$\neg p \vee \neg q$
F	F	T
F	T	T
T	F	T
T	T	F

Negating AND

p	q	$\neg(p \wedge q)$
F	F	T
F	T	T
T	F	T
T	T	F

p	q	$\neg p \vee \neg q$
F	F	T
F	T	T
T	F	T
T	T	F

These two statements
are always the same!

Logical Equivalence

- If two propositional logic statements φ and ψ always have the same truth values as one another, they are called **logically equivalent**.
- We denote this by $\varphi \equiv \psi$.
- \equiv is **not** a connective. Connectives are a part of logic statements; \equiv is something used to describe logic statements.
 - It is part of the **metalanguage** rather than the **language**.
- If $\varphi \equiv \psi$, we can modify any propositional logic formula containing φ by replacing it with ψ .
 - This is not true when we talk about first-order logic; we'll see why later.

De Morgan's Laws

- Using truth tables, we concluded that

$$\neg(p \wedge q) \equiv \neg p \vee \neg q$$

- We can also use truth tables to show that

$$\neg(p \vee q) \equiv \neg p \wedge \neg q$$

- These two equivalences are called **De Morgan's Laws**.

More Negations

- When is $p \rightarrow q$ false?
- **Answer:** p must be true and q must be false.
- In propositional logic:

$$p \wedge \neg q$$

- Is the following true?

$$\neg(p \rightarrow q) \equiv p \wedge \neg q$$

Negating Implications

Negating Implications

p	q	$\neg(p \rightarrow q)$
F	F	
F	T	
T	F	
T	T	

Negating Implications

p	q	$\neg(p \rightarrow q)$
F	F	T
F	T	T
T	F	F
T	T	T

Negating Implications

p	q	$\neg(p \rightarrow q)$	
F	F	F	T
F	T	F	T
T	F	T	F
T	T	F	T

Negating Implications

p	q	$\neg(p \rightarrow q)$	
F	F	F	T
F	T	F	T
T	F	T	F
T	T	F	T

Negating Implications

p	q	$\neg(p \rightarrow q)$		p	q	$p \wedge \neg q$
F	F	F	T	F	F	
F	T	F	T	F	T	
T	F	T	F	T	F	
T	T	F	T	T	T	

Negating Implications

p	q	$\neg(p \rightarrow q)$	
F	F	F	T
F	T	F	T
T	F	T	F
T	T	F	T

p	q	$p \wedge \neg q$
F	F	F
F	T	F
T	F	T
T	T	T

Negating Implications

p	q	$\neg(p \rightarrow q)$	
F	F	F	T
F	T	F	T
T	F	T	F
T	T	F	T

p	q	$p \wedge \neg q$	
F	F	F	T
F	T	F	F
T	F	T	T
T	T	T	F

Negating Implications

p	q	$\neg(p \rightarrow q)$	
F	F	F	T
F	T	F	T
T	F	T	F
T	T	F	T

p	q	$p \wedge \neg q$	
F	F	F	T
F	T	F	F
T	F	T	T
T	T	T	F

Negating Implications

p	q	$\neg(p \rightarrow q)$	
F	F	F	T
F	T	F	T
T	F	T	F
T	T	F	T

p	q	$p \wedge \neg q$	
F	F	F	T
F	T	F	F
T	F	T	T
T	T	F	F

Negating Implications

p	q	$\neg(p \rightarrow q)$		p	q	$p \wedge \neg q$	
F	F	F	T	F	F	F	T
F	T	F	T	F	T	F	F
T	F	T	F	T	F	T	T
T	T	F	T	T	T	F	F

$$\neg(p \rightarrow q) \equiv p \wedge \neg q$$

An Important Observation

- We have just proven that

$$\neg(p \rightarrow q) \equiv p \wedge \neg q$$

- If we negate both sides, we get that

$$p \rightarrow q \equiv \neg(p \wedge \neg q)$$

- By De Morgan's laws:

$$p \rightarrow q \equiv \neg(p \wedge \neg q)$$

$$p \rightarrow q \equiv \neg p \vee \neg\neg q$$

$$p \rightarrow q \equiv \neg p \vee q$$

- Thus $p \rightarrow q \equiv \neg p \vee q$

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$$p \rightarrow q \equiv \neg p \vee q$$

- Thus $p \rightarrow q \equiv \neg p \vee q$

If p is false, the whole thing is true and we gain no information. If p is true, then q has to be true for the whole expression to be true.

Another Idea

- We've just shown that $\neg(p \rightarrow q) \equiv p \wedge \neg q$.
- Is it also true that $\neg(p \rightarrow q) \equiv p \rightarrow \neg q$?
- Let's go check!

$\neg(p \rightarrow q)$ and $p \rightarrow \neg q$

$\neg(p \rightarrow q)$ and $p \rightarrow \neg q$

p	q	$\neg(p \rightarrow q)$
F	F	
F	T	
T	F	
T	T	

$\neg(p \rightarrow q)$ and $p \rightarrow \neg q$

p	q	$\neg(p \rightarrow q)$	
F	F	F	T
F	T	F	T
T	F	T	F
T	T	F	T

$\neg(p \rightarrow q)$ and $p \rightarrow \neg q$

p	q	$\neg(p \rightarrow q)$	
F	F	F	T
F	T	F	T
T	F	T	F
T	T	F	T

p	q	$p \rightarrow \neg q$
F	F	
F	T	
T	F	
T	T	

$\neg(p \rightarrow q)$ and $p \rightarrow \neg q$

p	q	$\neg(p \rightarrow q)$
F	F	F
F	T	F
T	F	T
T	T	F

p	q	$p \rightarrow \neg q$
F	F	F
F	T	F
T	F	T
T	T	T

$\neg(p \rightarrow q)$ and $p \rightarrow \neg q$

p	q	$\neg(p \rightarrow q)$	
F	F	F	T
F	T	F	T
T	F	T	F
T	T	F	T

p	q	$p \rightarrow \neg q$	
F	F	F	T
F	T	F	F
T	F	T	T
T	T	T	F

$\neg(p \rightarrow q)$ and $p \rightarrow \neg q$

p	q	$\neg(p \rightarrow q)$
F	F	F
F	T	F
T	F	T
T	T	F

p	q	$p \rightarrow \neg q$
F	F	T
F	T	F
T	F	T
T	T	F

$\neg(p \rightarrow q)$ and $p \rightarrow \neg q$

p	q	$\neg(p \rightarrow q)$
F	F	F
F	T	F
T	F	T
T	T	F

p	q	$p \rightarrow \neg q$
F	F	T
F	T	F
T	F	T
T	T	F

$\neg(p \rightarrow q)$ and $p \rightarrow \neg q$

p	q	$\neg(p \rightarrow q)$
F	F	F
F	T	F
T	F	T
T	T	F

p	q	$p \rightarrow \neg q$
F	F	T
F	T	F
T	F	T
T	T	F

$\neg(p \rightarrow q)$ and $p \rightarrow \neg q$

p	q	$\neg(p \rightarrow q)$
F	F	F
F	T	F
T	F	T
T	T	F

p	q	$p \rightarrow \neg q$
F	F	T
F	T	F
T	F	T
T	T	F

$\neg(p \rightarrow q)$ and $p \rightarrow \neg q$

p	q	$\neg(p \rightarrow q)$	
F	F	F	T
F	T	F	T
T	F	T	F
T	T	F	T

p	q	$p \rightarrow \neg q$	
F	F	F	T
F	T	F	F
T	F	T	T
T	T	T	F

$\neg(p \rightarrow q)$ and $p \rightarrow \neg q$

p	q	$\neg(p \rightarrow q)$		p	q	$p \rightarrow \neg q$	
F	F	F	T	F	F	F	T
F	T	F	T	F	T	F	F
T	F	T	F	T	F	T	T
T	T	F	T	T	T	T	F

These are not the same thing!

To prove that $p \rightarrow q$ is false, do **not** prove
 $p \rightarrow \neg q$.

Instead, prove that $p \wedge \neg q$ is true.

Analyzing Proof Techniques

Proof by Contrapositive

- Recall that to prove that $p \rightarrow q$, we can also show that $\neg q \rightarrow \neg p$.
- Let's verify that $p \rightarrow q \equiv \neg q \rightarrow \neg p$.

The Contrapositive

p	q	$p \rightarrow q$
F	F	
F	T	
T	F	
T	T	

The Contrapositive

p	q	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T

The Contrapositive

p	q	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T

p	q	$\neg q \rightarrow \neg p$
F	F	
F	T	
T	F	
T	T	

The Contrapositive

p	q	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T

p	q	$\neg q \rightarrow \neg p$
F	F	T
F	T	F
T	F	T
T	T	F

The Contrapositive

p	q	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T

p	q	$\neg q \rightarrow \neg p$
F	F	T
F	T	F
T	F	T
T	T	F

The Contrapositive

p	q	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T

p	q	$\neg q \rightarrow \neg p$
F	F	T
F	T	F
T	F	T
T	T	F

The Contrapositive

p	q	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T

p	q	$\neg q \rightarrow \neg p$
F	F	T
F	T	F
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T	T	F

The Contrapositive

p	q	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T

p	q	$\neg q \rightarrow \neg p$
F	F	T
F	T	F
T	F	T
T	T	F

The Contrapositive

p	q	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T

p	q	$\neg q \rightarrow \neg p$
F	F	T
F	T	F
T	F	T
T	T	F

The Contrapositive

p	q	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T

p	q	$\neg q$	$\neg q \rightarrow \neg p$	$\neg p$
F	F	T	T	T
F	T	F	T	T
T	F	T	F	F
T	T	F	T	F

The Contrapositive

p	q	$p \rightarrow q$	p	q	$\neg q \rightarrow \neg p$
F	F	T	F	F	T
F	T	T	F	T	T
T	F	F	T	F	F
T	T	T	T	T	T

$$p \rightarrow q \equiv \neg q \rightarrow \neg p$$

Why All This Matters

- Suppose we want to prove the following statement:

“If $x + y = 16$, then $x \geq 8$ or $y \geq 8$ ”

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$$\neg(x \geq 8 \vee y \geq 8) \rightarrow \neg(x + y = 16)$$

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$$x < 8 \wedge y < 8 \rightarrow x + y \neq 16$$

“If $x < 8$ and $y < 8$, then $x + y \neq 16$ ”

Theorem: If $x + y = 16$, then either $x \geq 8$ or $y \geq 8$.

Proof: By contrapositive. We prove that if $x < 8$ and $y < 8$, then $x + y \neq 16$. To see this, note that

$$\begin{aligned}x + y &< 8 + y \\ &< 8 + 8 \\ &= 16\end{aligned}$$

So $x + y < 16$, so $x + y \neq 16$. ■

Why This Matters

- Propositional logic is a tool for reasoning about how various statements affect one another.
- To better understand how to prove a result, it often helps to translate what you're trying to prove into propositional logic first.
- **Note:** To truly reason about proofs, we need the more expressive power of **first-order logic**, which we'll talk about next time.

Proof by Contradiction

- The general structure of a proof by contradiction is
 - To show p , assume p is false.
 - Show that p being false implies something that cannot be true.
 - Conclude, therefore, that p is true.
- What does this look like in propositional logic?

Proof by Contradiction

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- What does this look like in propositional logic?

$$\neg p$$

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$$\neg p \rightarrow \perp$$

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 - To show p , assume p is false.
 - Show that p being false implies something that cannot be true.
 - Conclude, therefore, that p is true.
- What does this look like in propositional logic?

$$(\neg p \rightarrow \perp) \rightarrow p$$

Proof by Contradiction

Proof by Contradiction

p	$(\neg p \rightarrow \perp) \rightarrow p$
F	
T	

Proof by Contradiction

p	$(\neg p \rightarrow \perp) \rightarrow p$
F	T
T	F

Proof by Contradiction

p	$(\neg p \rightarrow \perp) \rightarrow p$
F	T
T	F

Proof by Contradiction

p	$(\neg p \rightarrow \perp) \rightarrow p$		
F	T	F	F
T	F		F

Proof by Contradiction

p	$(\neg p \rightarrow \perp) \rightarrow p$		
F	T	F	F
T	F	T	F

Proof by Contradiction

p	$(\neg p \rightarrow \perp) \rightarrow p$			p
F	T	F	F	F
T	F	T	F	T

Proof by Contradiction

p	$(\neg p \rightarrow \perp) \rightarrow p$				
F	T	F	F	T	F
T	F	T	F		T

Proof by Contradiction

p	$(\neg p \rightarrow \perp) \rightarrow p$				
F	T	F	F	T	F
T	F	T	F	T	T

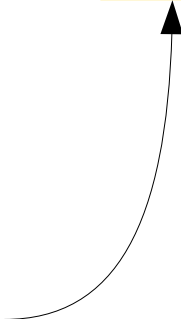
Proof by Contradiction

p	$(\neg p \rightarrow \perp)$			\rightarrow	p
F	T	F	F	T	F
T	F	T	F	T	T

Proof by Contradiction

p	$(\neg p \rightarrow \perp)$			\rightarrow	p
F	T	F	F	T	F
T	F	T	F	T	T

This statement
is always true!



Tautologies

- A **tautology** is a statement that is always true.
- Examples:
 - \top
 - $p \vee \neg p$ (the **Law of the Excluded Middle**)
 - $\perp \rightarrow p$ (**vacuous truth**)
- Once a tautology has been proven, we can use that tautology anywhere.

Next Time

- **First-Order Logic**
 - How do we reason about multiple objects and their properties?