CS 106B, Lecture 8 Recursion

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Plan for Today

- Learn a powerful algorithmic technique called *recursion*
 - Exploit self-similarity in problems
- We will spend several days on recursion don't worry if it doesn't make sense today
 - Goal: do as many examples as we can
 - You should practice: <u>CodeStepByStep</u>, section problems, or examples from the textbook
 - Highly encourage the reading for this week!

Recursion

- recursion: Function that calls itself
 - Solving a problem using recursion depends on solving smaller (simpler) occurrences of the same problem until the problem is simple enough that you can solve it directly
 - Key question: "How is this problem self-similar?" what are the smaller subproblems that make up the bigger problem?
- Occurs in many places in code and in real world:
 - Looking up a word in dictionary may involve looking up other words
 - Nested structures (trees, file folders, collections) can be self-similar

Recursive Programming

- **recursive programming**: Writing <u>functions that call themselves</u> to solve problems that are recursive in nature.
 - An equally powerful substitute for *iteration* (loops)
 - Particularly well-suited to solving certain types of problems
 - Leads to **elegant**, simplistic, short code (when used well)
 - A key component of many of our assignments in this course

Non-recursive factorial

```
// Returns n!, or 1 * 2 * 3 * 4 * ... * n.
// Assumes n >= 1.
int factorial(int n) {
    int total = 1;
    for (int i = 1; i <= n; i++) {
        total *= i;
      }
    return total;
}</pre>
```

• Important observations:

$$0! = 1! = 1$$

$$4! = 4 * 3 * 2 * 1$$

$$5! = 5 * 4 * 3 * 2 * 1$$

$$= 5 * 4!$$

Recursive factorial

- The recursive code handles a small part of the overall task (multiplying by *n*), then makes a recursive call to handle the rest.
 - The recursive version is written without using any loops.
 - Recursion *replaces* the loop
 - We separate the code into a *base case* (a simple case that does not make any recursive calls), and a *recursive case*.

Recursive stack trace

```
int factorial(int n) { // 4
    if (n <= 1) {
                                       // base case
        return 1;
    } else {
       return n * factorial(n - 1); // recursive case
   int factorial(int n) { // 3
       if (n <= 1) {
                                          // base case
           return 1;
       } else {
           return n * factorial(n - 1); // recursive case
      int factorial(int n) { // 2
   }
          if (n <= 1) {
                                              // base case
              return 1;
           } else {
              return n * factorial(n - 1); // recursive case
         int factorial(int n) { // 1
      }
              if (n <= 1) {
                                                 // base case
                  return 1;
              } else {
                  return n * factorial(n - 1); // recursive case
              }
```

Recursion and cases

- Every recursive algorithm involves at least 2 cases:
 - **base case**: A simple occurrence that can be answered directly
 - recursive case: A more complex occurrence of the problem that cannot be directly answered, but can instead be described in terms of smaller occurrences of the same problem
 - Key idea: In a recursive piece of code, you handle a small part of the overall task yourself (usually the work involves modifying the results of the smaller problems), then make a recursive call to handle the rest.
 - Ask yourself, "How is this task self-similar?"
 - "How can I describe this algorithm in terms of a smaller or simpler version of itself?"

Three Rules of Recursion

- Every (valid) input must have a case (either recursive or base)
- There **must** be a base case that makes no recursive calls
- The recursive case must make the problem simpler and make forward progress to the base case

Recursive tracing

• Consider the following recursive function:

```
int mystery(int n) {
    if (n < 10) {
        return n;
    } else {
        int a = n / 10;
        int b = n % 10;
        return mystery(a + b);
    }
}</pre>
```

Q: What is the result of: mystery(648) ?
A. 8
B. 9
C. 54
D. 72
E. 648

Recursive stack trace



isPalindrome exercise

• Write a recursive function isPalindrome that accepts a string and returns true if it reads the same forwards as backwards.

\rightarrow true
\rightarrow true
\rightarrow true
\rightarrow true
\rightarrow true
\rightarrow false
\rightarrow false
\rightarrow false
\rightarrow false

- What is a good **base case**?

isPalindrome

- What is our stopping point (*base case*)?
- How is this problem *self-similar*?
- What is the minimum *amount of work*?
- How can we make the problem *simpler* by doing the least amount of work?

isPalindrome

- What is our stopping point (*base case*)?
 - Empty string or string of length 1
- How is this problem *self-similar*?
 - Palindromes can be written as: x[SMALLER_PALINDROME]x, where x stands for some letter
- What is the minimum *amount of work*?
 - Testing the equality of outside characters
- How can we make the problem *simpler* by doing the least amount of work?
 - Peel off the outside characters and test if the middle is a palindrome

isPalindrome solution

```
// Returns true if the given string reads the same
// forwards as backwards.
// Trivially true for empty or 1-letter strings.
bool isPalindrome(string s) {
    if (s.length() < 2) { // base case</pre>
        return true;
    } else {
                        // recursive case
        if (s[0] != s[s.length() - 1]) {
            return false;
        }
        string middle = s.substr(1, s.length() - 2);
        return isPalindrome(middle);
    }
```

Announcements

- Homework 2 due on Wednesday at **5PM**
- Homework 1 grades will be released by your section leader on or before Wednesday
- Alternate midterms are being scheduled this week. Keep an eye out for an email from Kate

Multiple calls tracing

```
int mystery(int n) {
    if (n < 10) {
        return (10 * n) + n;
    } else {
        int a = mystery(n / 10);
        int b = mystery(n % 10);
        return (100 * a) + b;
    }
}</pre>
```

Q: What is the result of: mystery(348) ?
 A. 3828 B. 348348 C. 334488 D. 80403 E. none

Multiple calls tracing



Recursive Big O

- Below is the "pseudocode" for finding Big O of a function
 - Note that this is not real code; this is to show the recursive nature of finding Big O
 - Self-similarity: find Big O of smaller code blocks and combine them
 - This Big O pseudocode doesn't cover function calls and some other cases (for pedagogical purposes)

findBigO(codeSnippet):

- if codeSnippet is a single statement:
 return O(1)
- if codeSnippet is loop:

return number of times loop runs * findBigO(loop inside)
for codeBlock in codeSnippet:
 return the sum of findBigO(codeBlock)

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```
for (int i = 0; i < N * N; i += 3) {
    for (int j = 3; j <= 219; j++) {
        cout << "sum: " << i + j << endl;
    }
}</pre>
```

```
cout << "Have a nice Life!" << endl;</pre>
```

findBigO(codeSnippet):

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for (int i = 0; i < N * N; i += 3) {
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        cout << "sum: " << i + j << endl;
    }
}</pre>
```

findBigO(codeSnippet):

```
if codeSnippet is a single statement:
       return O(1)
                                       O(N^2)
   if codeSnippet is loop:
       return number of times loop runs * findBigO(loop inside)
   for codeBlock in codeSnippet:
       return the sum of findBigO(codeBlock)
for (int i = 0; i < N * N; i += 3) {
    for (int j = 3; j <= 219; j++) {
        cout << "sum: " << i + j << endl;</pre>
    }
```

findBigO(codeSnippet):

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for (int i = 0; i < N * N; i += 3) {</pre>
    for (int j = 3; j <= 219; j++) {
        cout << "sum: " << i + j << endl; _</pre>
    }
```

```
cout << "Have a nice Life!" << endl;</pre>
```

```
findBigO(codeSnippet):
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    return O(1)
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```
for (int i = 0; i < N * N; i += 3) {
    for (int j = 3; j <= 219; j++) {
        cout << "sum: " << i + j << endl;
    }
}</pre>
```

```
cout << "Have a nice Life!" << endl;</pre>
```

findBigO(codeSnippet):

if codeSnippet is a single statement: return 0(1) if codeSnippet is loop: _____O(1) _____O(1) return number of times loop runs * findBigO(loop inside) for codeBlock in codeSnippet: return the sum of findBigO(codeBlock) for (int i = 0; i < N * N; i += 3) {</pre>

}

findBigO(codeSnippet):

if codeSnippet is a single statement: return O(1) $O(N^2)$ if codeSnippet is loop: return number of times loop runs * findBigO(oop instae) for codeBlock in codeSnippet: return the sum of findBigO(codeBlock) for (int i = 0; i < N * N; i += 3) { for (int j = 3; j <= 219; j++) {</pre> cout << "sum: " << i + j << endl;</pre> }

```
findBigO(codeSnippet):
   if codeSnippet is a single statement:
       return O(1)
   if codeSnippet is loop:
       return number of times loop runs * findBigO(loop inside)
   for codeBlock in codeSnippet:
       return the sum of findBigO(codeBlock) O(N<sup>2</sup>) +
for (int i = 0; i < N * N; i += 3) {</pre>
    for (int j = 3; j <= 219; j++) {
        cout << "sum: " << i + j << endl;</pre>
    }
```

```
findBigO(codeSnippet):
```

```
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if codeSnippet is loop:
    return number of times loop runs * findBigO(loop inside)
for codeBlock in codeSnippet:
    return the sum of findBigO(codeBlock)
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```
for (int i = 0; i < N * N; i += 3) {
    for (int j = 3; j <= 219; j++) {
        cout << "sum: " << i + j << endl;
    }
}</pre>
```

```
findBigO(codeSnippet):
   if codeSnippet is a single statement:
       return O(1)
   if codeSnippet is loop:
       return number of times loop runs * findBigO(loop inside)
   for codeBlock in codeSnippet:
       return the sum of findBigO(codeBlock) O(N<sup>2</sup>) + O(1)
for (int i = 0; i < N * N; i += 3) {
    for (int j = 3; j <= 219; j++) {
        cout << "sum: " << i + j << endl;</pre>
    }
```

```
cout << "Have a nice Life!" << endl;</pre>
```

findBigO(codeSnippet):

- if codeSnippet is a single statement:
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return number of times loop runs * findBigO(loop inside)
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final result: O(N²)

```
for (int i = 0; i < N * N; i += 3) {
   for (int j = 3; j <= 219; j++) {
      cout << "sum: " << i + j << endl;
   }
}</pre>
```

power exercise

- Write a function **power** that accepts integer parameters for a base and exponent and computes base ^ exponent.
 - Write a <u>recursive</u> version of this function (one that calls itself).
 - Solve the problem without using any loops.
 - What is our stopping point (base case)?
 - How is this problem *self-similar*?
 - What is the minimum amount of work?
 - How can we make the problem *simpler* by doing the least amount of work?

Initial solution

```
// Returns base ^ exp.
// Assumes exp >= 1.
int power(int base, int exp) {
    if (exp == 1) {
        return base;
    } else {
        return base * power(base, exp - 1);
    }
}
```

The call stack

• Each previous call waits for the next call to finish.

```
- cout << power(5, 3) << endl;</pre>
```



"Recursion Zen"

• The real, even simpler, base case is an exp of 0, not 1:

```
int power(int base, int exp) {
    if (exp == 0) {
        // base case; base^0 = 1
        return 1;
    } else {
        // recursive case: x^y = x * x^(y-1)
        return base * power(base, exp - 1);
    }
}
```

 Recursion Zen: The art of properly identifying the best set of cases for a recursive algorithm and expressing them elegantly.
 Opposite is arms-length recursion (our informal term)

Preconditions

- precondition: Something your code assumes is true when called.
 - Often documented as a comment on the function's header:

```
// Returns base ^ exp.
// Precondition: exp >= 0
int power(int base, int exp) {
```

- Stating a precondition doesn't really "solve" the problem, but it at least documents our decision and warns the client what not to do.
- What if the caller doesn't listen and passes a negative power anyway?
 What if we want to actually *enforce* the precondition?

Throwing exceptions

error(expression);

- In Stanford C++ lib's "error.h"
- Generates an exception that will crash the program, unless it has code to handle ("catch") the exception.
- alternative: throw something
 - *something* can be an int, a string, etc.
- Why would anyone ever *want* a program to crash?

power solution 2

```
// Returns base ^ exp.
// Precondition: exp >= 0
int power(int base, int exp) {
    if (exp < 0) {
        throw "illegal negative exponent";
    } else ...
    ...</pre>
```

}