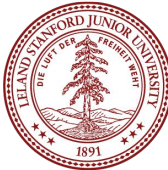


Introduction to Recursion

What's been the most challenging part of
Assignment 2 for you so far?
(put your answers the chat)



Roadmap

C++ basics

User/client

vectors + grids

stacks + queues

sets + maps

Core
Tools

testing

algorithmic
analysis

recursive
problem-solving

Object-Oriented
Programming

Implementation

arrays

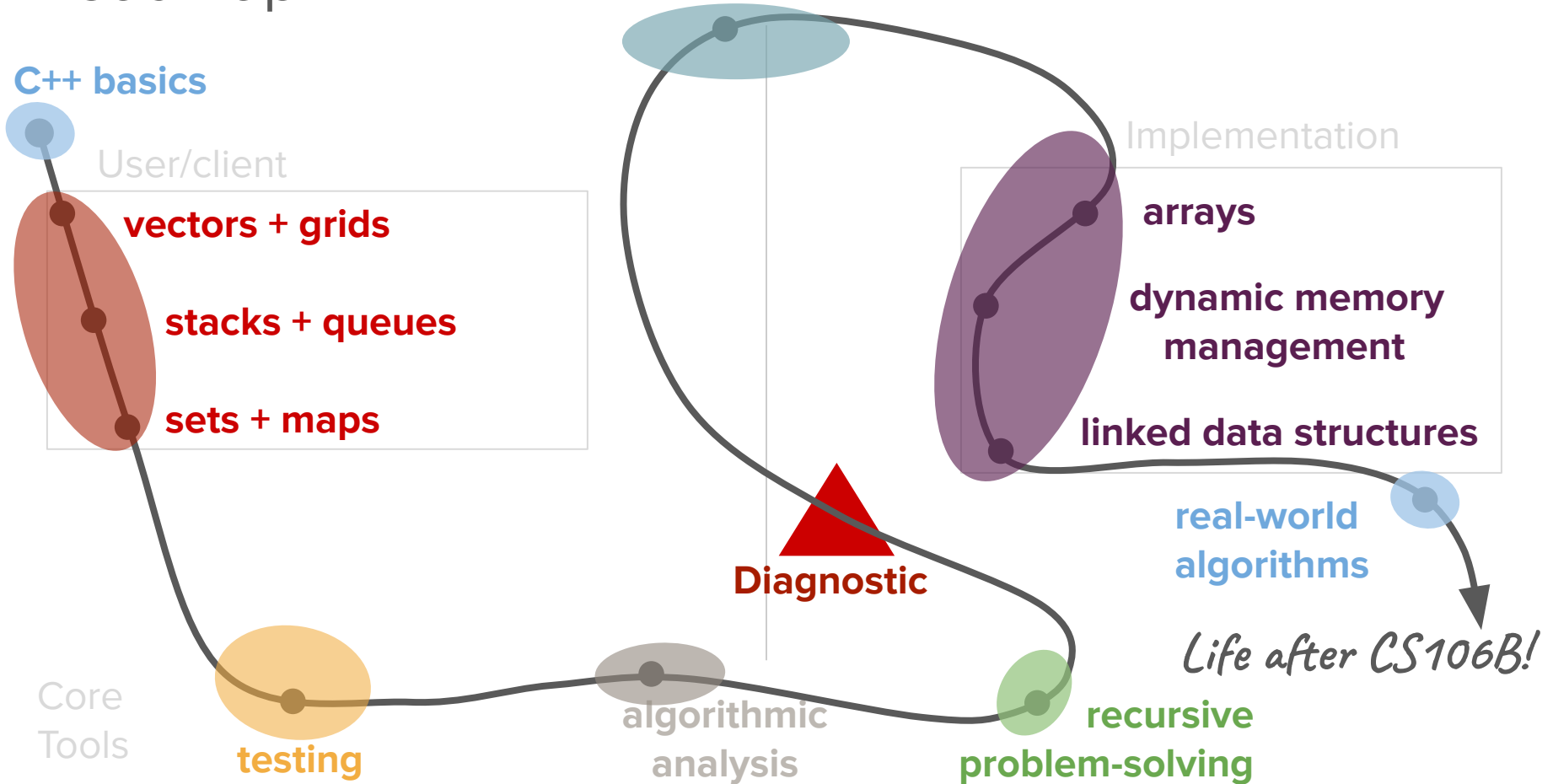
dynamic memory
management

linked data structures

real-world
algorithms

Life after CS106B!

Diagnostic



Roadmap

C++ basics

User/client

vectors + grids

stacks + queues

sets + maps

Core
Tools

testing

algorithmic
analysis

recursive
problem-solving

Object-Oriented
Programming

Implementation

arrays

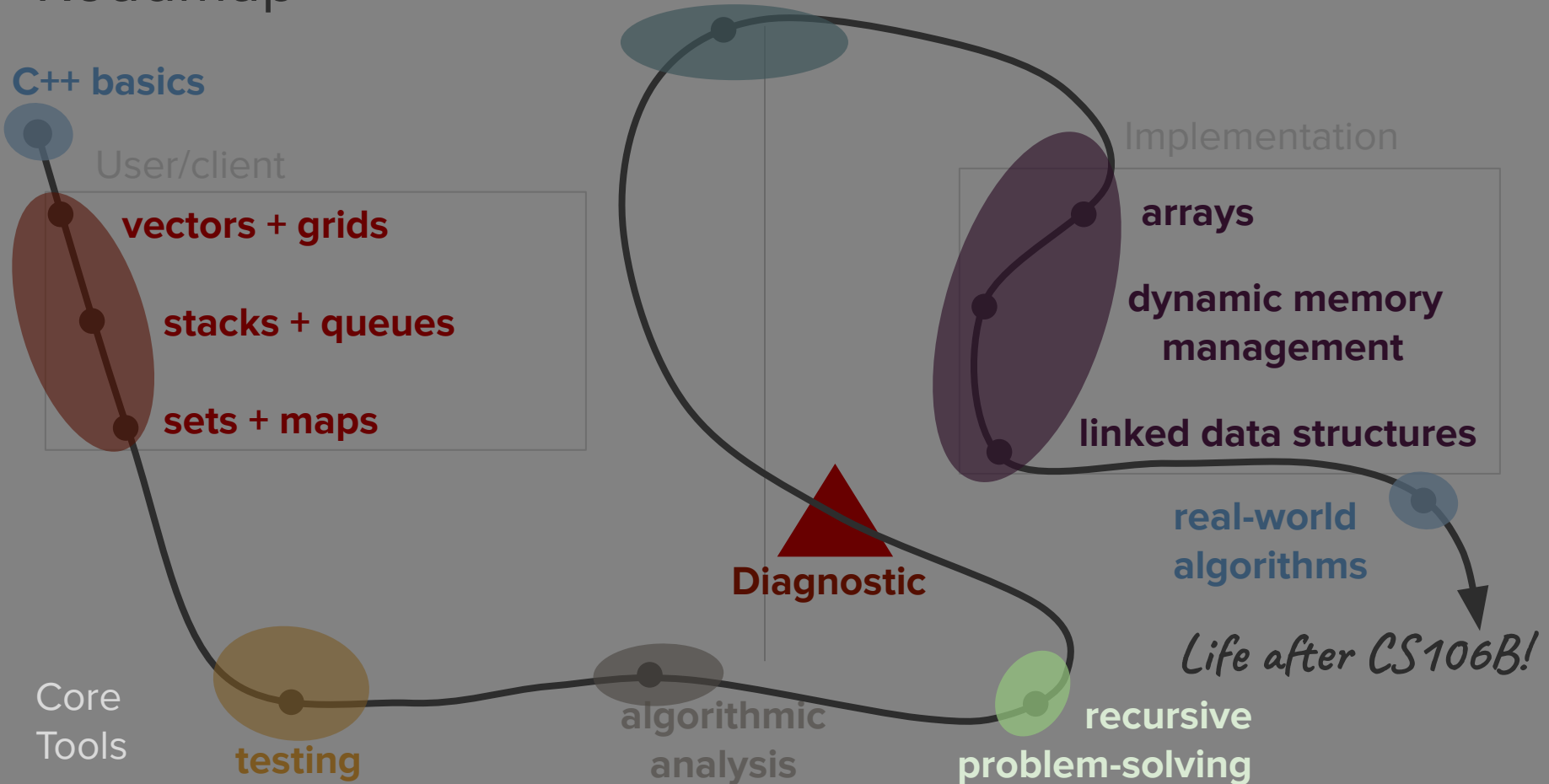
dynamic memory
management

linked data structures

real-world
algorithms

Life after CS106B!

Diagnostic



Today's question

How can we take
advantage of self-similarity
within a problem to solve it
more elegantly?

Today's topics

1. Review
2. Defining recursion
3. Recursion + Stack Frames
(e.g. factorials)
4. Recursive Problem-Solving
(e.g. string reversal)

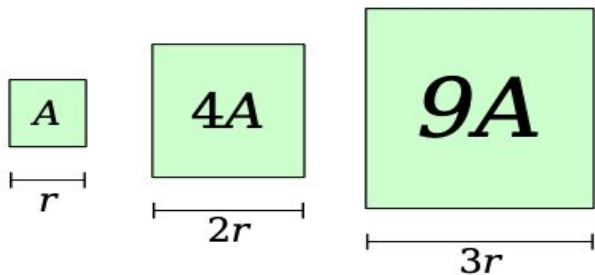
Review

(Big O)

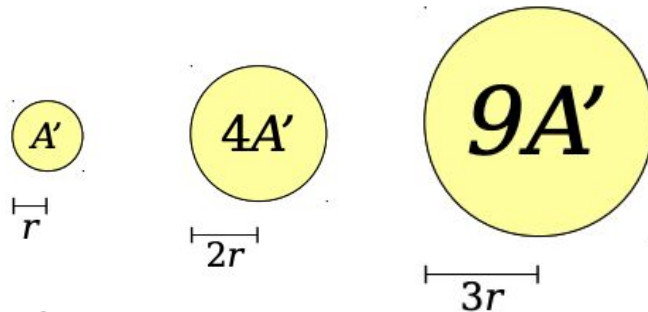
Big-O Notation

- **Big-O notation** is a way of quantifying the rate at which some quantity grows.
- Example:
 - A square of side length r has area $O(r^2)$.
 - A circle of radius r has area $O(r^2)$.

This just says that these quantities grow at the same relative rates. It does not say that they're equal!



*Doubling r increases area 4x
Tripling r increases area 9x*



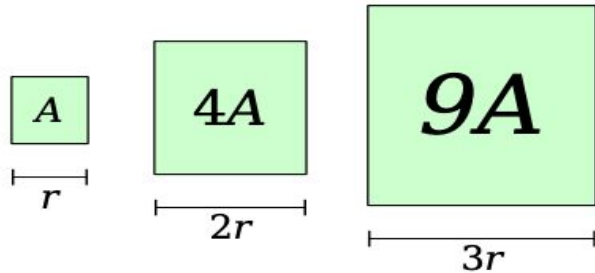
*Doubling r increases area 4x
Tripling r increases area 9x*

Big-O Notation

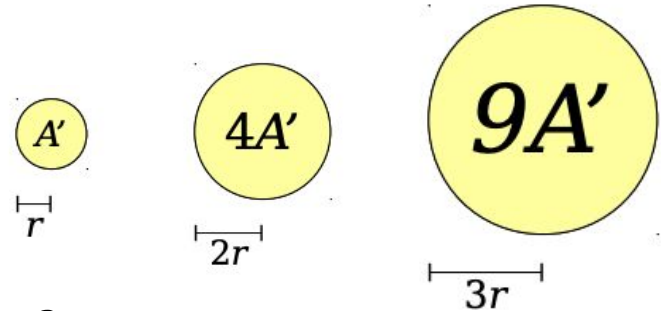
With respect to a given input variable!



- **Big-O notation** is a way of quantifying the **rate at which some quantity grows**.
- Example:
 - A square of side length r has area $O(r^2)$.
 - A circle of radius r has area $O(r^2)$.



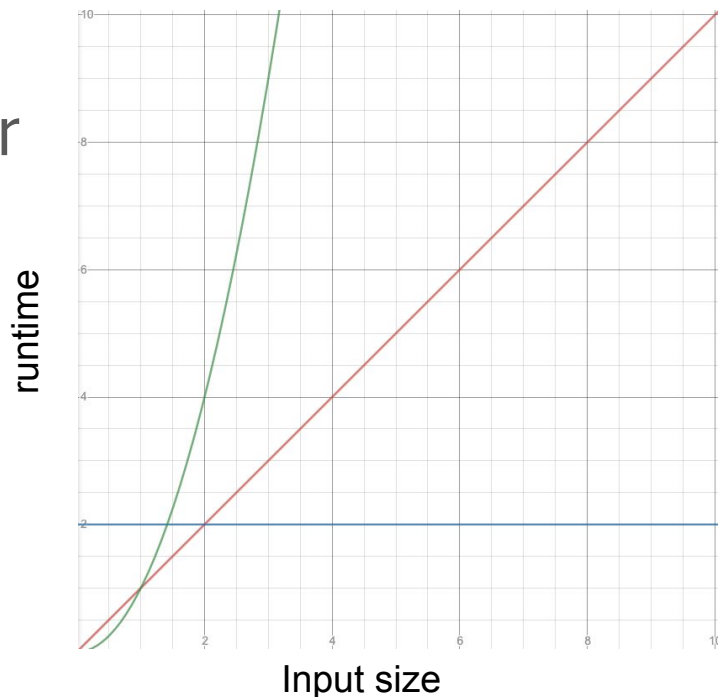
*Doubling r increases area 4x
Tripling r increases area 9x*



*Doubling r increases area 4x
Tripling r increases area 9x*

Efficiency Categorizations So Far

- Constant Time – $O(1)$
 - Super fast, this is the best we can hope for!
 - Euclid's Algorithm for Perfect Numbers
- Linear Time – $O(n)$
 - This is okay; we can live with this
- Quadratic Time – $O(n^2)$
 - This can start to slow down really quickly
 - Exhaustive Search for Perfect Numbers
- How do all the ADT operations we've seen so far fall into these categories?



ADT Big-O Matrix

● Vectors

- `.size()` - $O(1)$
- `.add()` - $O(1)$
- `v[i]` - $O(1)$
- `.insert()` - $O(n)$
- `.remove()` - $O(n)$
- `.clear()` - $O(n)$
- `traversal` - $O(n)$ }

● Grids

- `.numRows()` / `.numCols()`
- $O(1)$
- `g[i][j]` - $O(1)$
- `.inBounds()` - $O(1)$
- `traversal` - $O(n^2)$

● Queues

- `.size()` - $O(1)$
- `.peek()` - $O(1)$
- `.enqueue()` - $O(1)$
- `.dequeue()` - $O(1)$
- `.isEmpty()` - $O(1)$
- `traversal` - $O(n)$

● Stacks

- `.size()` - $O(1)$
- `.peek()` - $O(1)$
- `.push()` - $O(1)$
- `.pop()` - $O(1)$
- `.isEmpty()` - $O(1)$
- `traversal` - $O(n)$

● Sets

- `.size()` - $O(1)$
- `.isEmpty()` - $O(1)$
- `.add()` - ???
- `.remove()` - ???
- `.contains()` - ???
- `traversal` - $O(n)$

● Maps

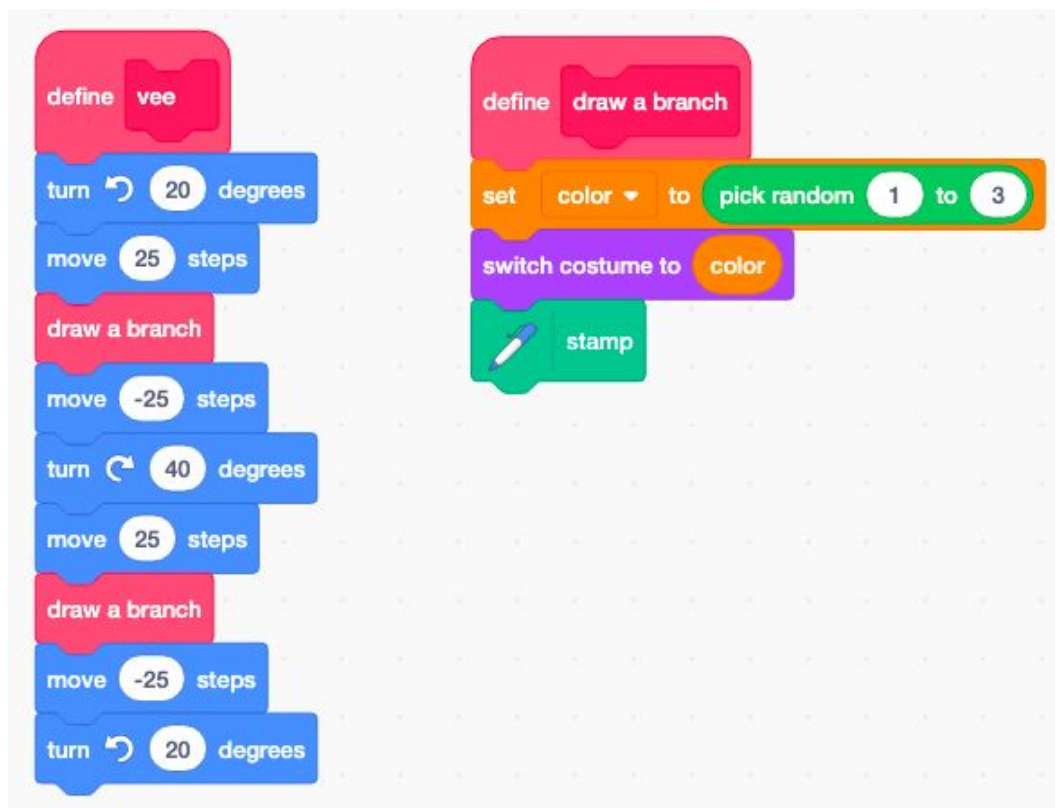
- `.size()` - $O(1)$
- `.isEmpty()` - $O(1)$
- `m[key]` - ???
- `.contains()` - ???
- `traversal` - $O(n)$

What is recursion?

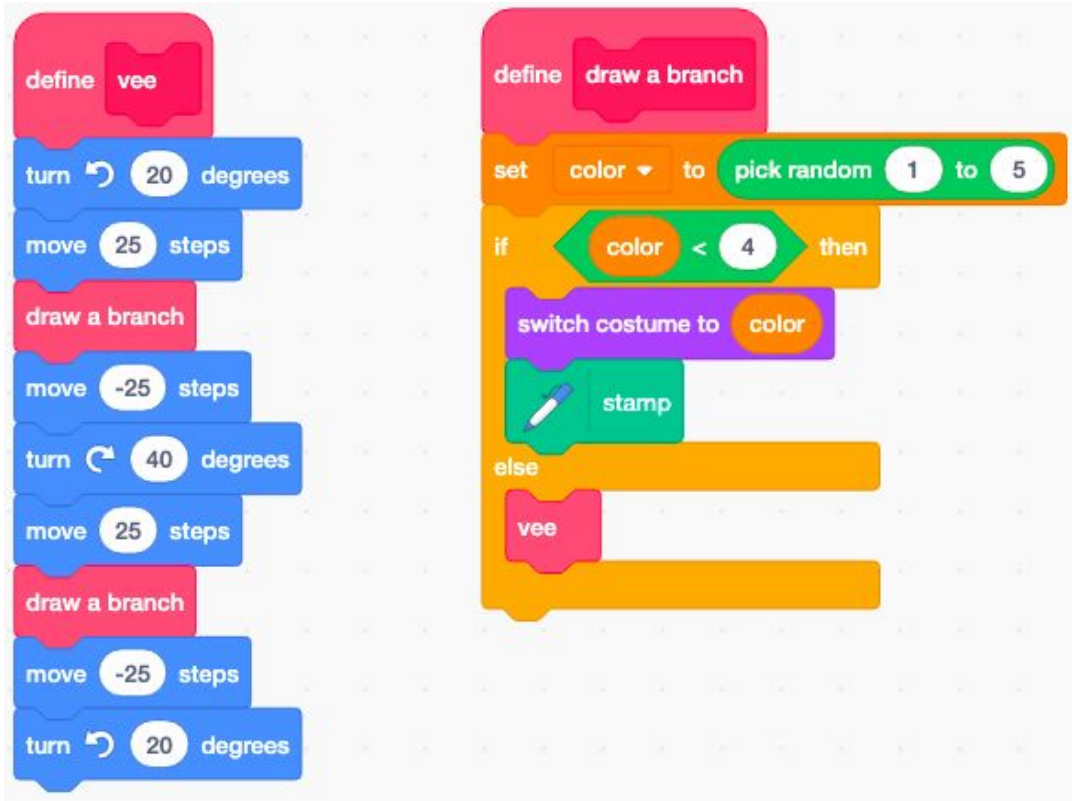
Activity: Vee

(<https://scratch.mit.edu/projects/409796637/>)

This code creates a “vee” shape with random colors.



Discuss in breakout rooms: What will this code do?



Discuss in breakout rooms: What will this code do?

The image shows two Scratch code blocks. The first block, 'define vee', consists of: a pink 'define vee' block, a blue 'turn 20 degrees' block, a blue 'move 25 steps' block, a pink 'draw a branch' block, a blue 'move -25 steps' block, a blue 'turn 40 degrees' block, a blue 'move 25 steps' block, a pink 'draw a branch' block, a blue 'move -25 steps' block, and a blue 'turn 20 degrees' block. The second block, 'define draw a branch', consists of: a pink 'define draw a branch' block, an orange 'set color to pick random 1 to 5' block (with '5' circled), an orange 'if color < 4 then' block containing a purple 'switch costume to color' block and a green 'stamp' block, and an orange 'else' block containing a pink 'vee' block.

Notice the differences

Demo: Recursive Vee

(<https://scratch.mit.edu/projects/409785610/>)

What is recursion?

Wikipedia: “Recursion occurs when a thing is defined in terms of itself.”



recursion



 All

 Books

 Images

 Videos

 News

 More

Settings

Tools

About 33,900,000 results (0.53 seconds)

Did you mean: ***recursion***

Definition

recursion

A problem-solving technique in which tasks are completed by reducing them into repeated, smaller tasks of the same form.

What is recursion?

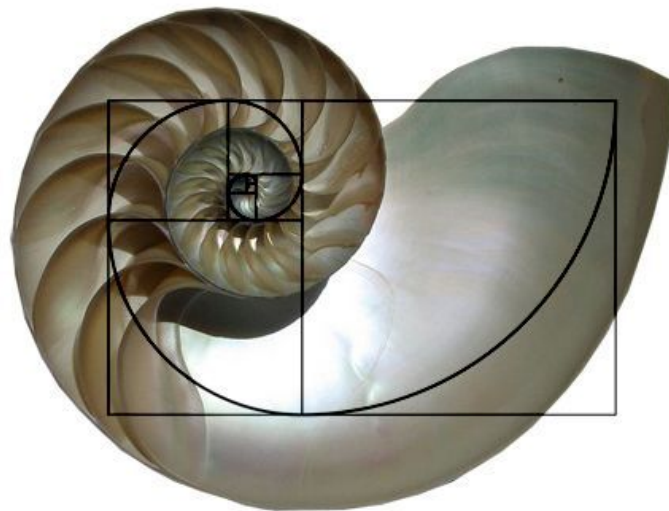
- A powerful substitute for iteration (loops)
 - We'll start off with seeing the difference between iterative vs. recursive solutions
 - Later in the week we'll see problems/tasks that can only be solved using recursion

What is recursion?

- A powerful substitute for iteration (loops)
 - We'll start off with seeing the difference between iterative vs. recursive solutions
 - Later in the week we'll see problems/tasks that can only be solved using recursion
- Results in elegant, often shorter code when used well

What is recursion?

- A powerful substitute for iteration (loops)
 - We'll start off with seeing the difference solutions
 - Later in the week we'll see problems/ta recursion
- Results in elegant, often shorter code when
- Often applied to sorting and searching problems and can be used to express patterns seen in nature



What is recursion?

- A powerful substitute for iteration (loops)
 - We'll start off with seeing the difference between iterative vs. recursive solutions
 - Later in the week we'll see problems/tasks that can only be solved using recursion
- Results in elegant, often shorter code when used well
- Often applied to sorting and searching problems and can be used to express patterns seen in nature
- Will be part of many of our future assignments!

How many students
are in a lecture hall?

A [non-COVID] analogy

How many students are in the lecture hall?

- Let's suppose I want to find out how many people are at lecture today, but I don't want to walk around and count each person.
- I want to recruit your help, but I also want to minimize each individual's amount of work.

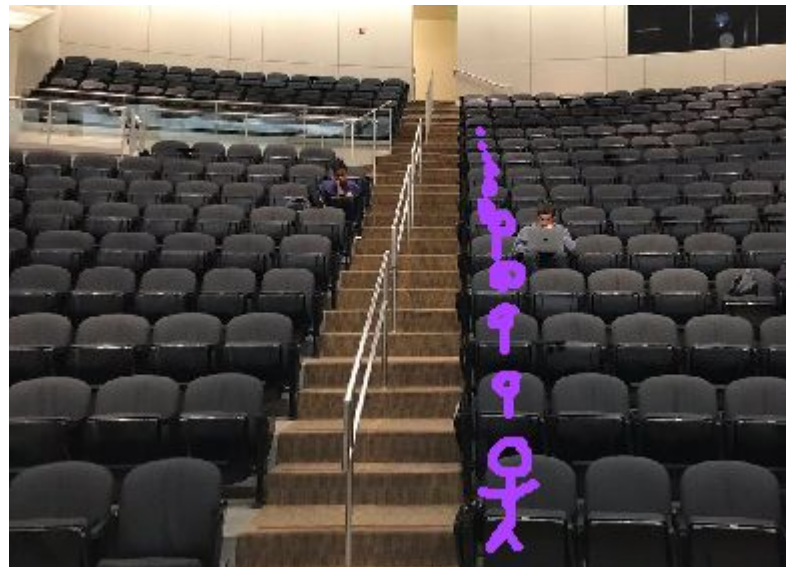
How many students are in the lecture hall?

- Let's suppose I want to find out how many people are at lecture today, but I don't want to walk around and count each person.
- I want to recruit your help, but I also want to minimize each individual's amount of work.

We can solve this problem recursively!

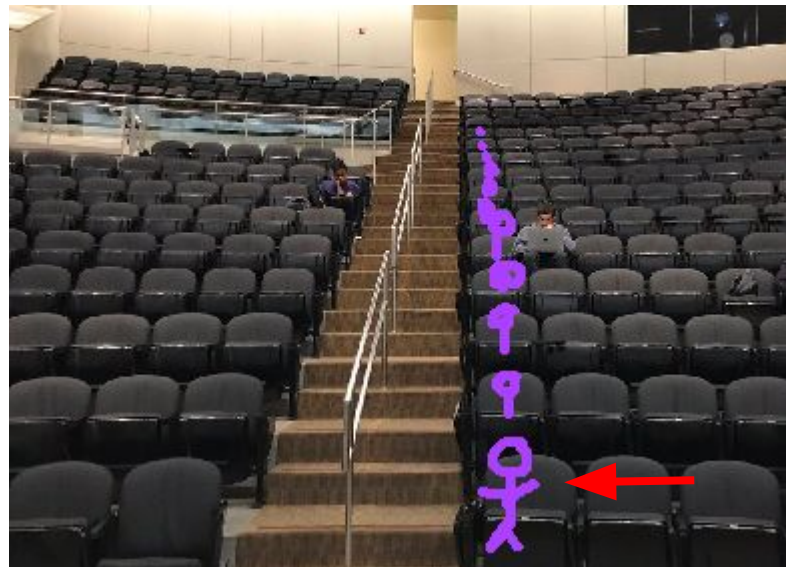
How many students are in the lecture hall?

- We'll focus on solving the problem for single "column" of students.



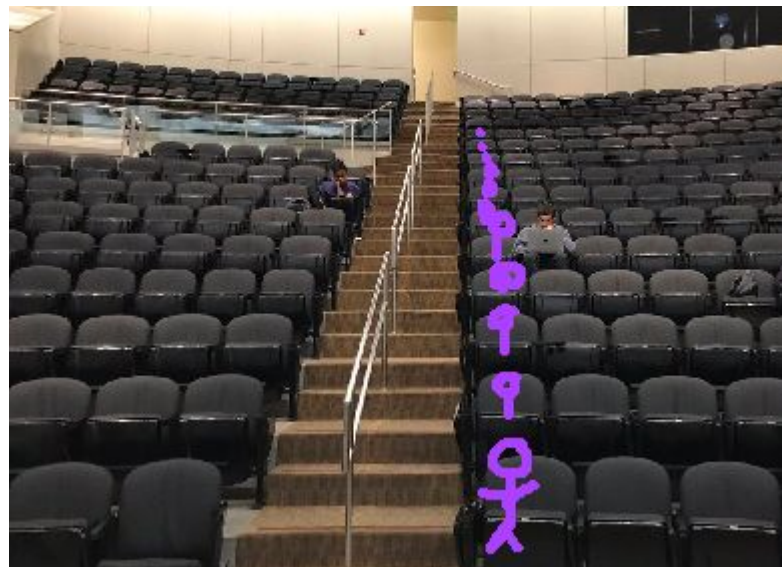
How many students are in the lecture hall?

- We'll focus on solving the problem for single “column” of students.
 - I go to the first person in the front row and ask: “How many people are sitting directly behind you in your ‘column’?”



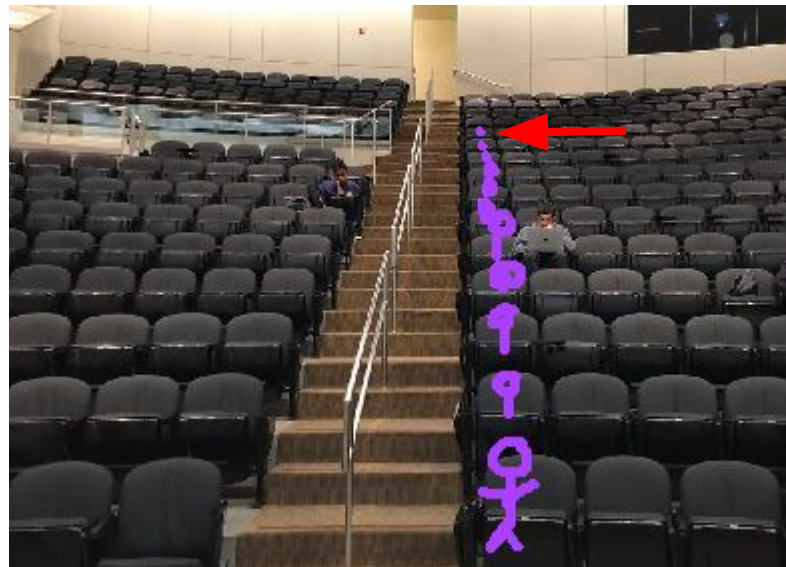
How many students are in the lecture hall?

- We'll focus on solving the problem for single "column" of students.
 - I go to the first person in the front row and ask: "How many people are sitting directly behind you in your 'column'?"
 - Student's algorithm:
 - If there is no one behind me, answer 0.
 - If someone is sitting behind me:
 - Ask that person: How many people are sitting directly behind you in your "column"?
 - When they respond with a value N , respond $(N + 1)$ to the person who asked me.



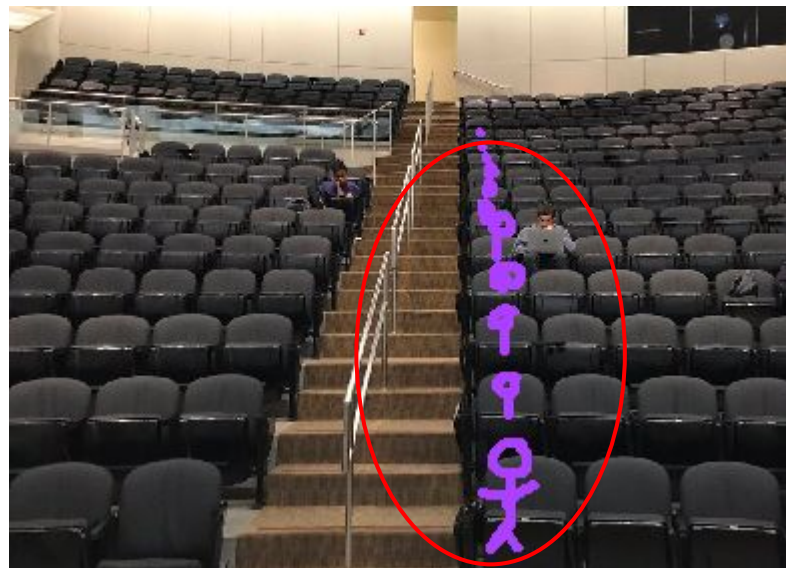
How many students are in the lecture hall?

- We'll focus on solving the problem for single "column" of students.
 - I go to the first person in the front row and ask: "How many people are sitting directly behind you in your 'column'?"
 - Student's algorithm:
 - **If there is no one behind me, answer 0.**
 - If someone is sitting behind me:
 - Ask that person: How many people are sitting directly behind you in your "column"?
 - When they respond with a value N , respond $(N + 1)$ to the person who asked me.



How many students are in the lecture hall?

- We'll focus on solving the problem for single "column" of students.
 - I go to the first person in the front row and ask: "How many people are sitting directly behind you in your 'column'?"
 - Student's algorithm:
 - If there is no one behind me, answer 0.
 - **If someone is sitting behind me:**
 - Ask that person: How many people are sitting directly behind you in your "column"?
 - When they respond with a value N , respond $(N + 1)$ to the person who asked me.



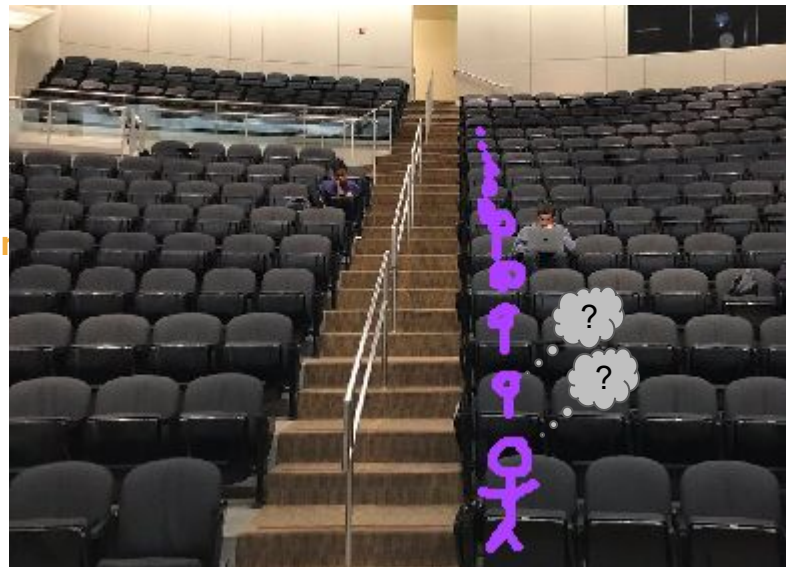
How many students are in the lecture hall?

- We'll focus on solving the problem for single "column" of students.
 - I go to the first person in the front row and ask: "How many people are sitting directly behind you in your 'column'?"
 - Student's algorithm:
 - If there is no one behind me, answer 0.
 - If someone is sitting behind me:
 - Ask that person: How many people are sitting directly behind you in your "column"?
 - When they respond with a value N , respond $(N + 1)$ to the person who asked me.



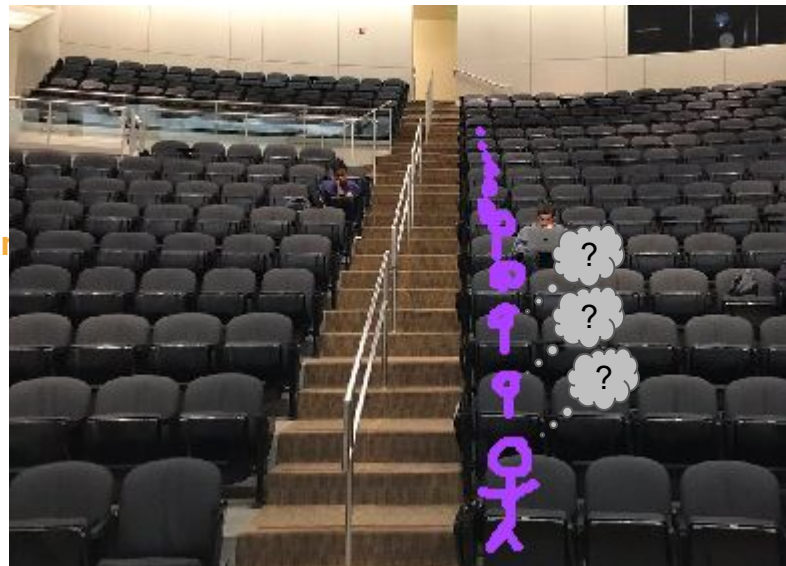
How many students are in the lecture hall?

- We'll focus on solving the problem for single "column" of students.
 - I go to the first person in the front row and ask: "How many people are sitting directly behind you in your 'column'?"
 - Student's algorithm:
 - If there is no one behind me, answer 0.
 - **If someone is sitting behind me:**
 - **Ask that person: How many people are sitting directly behind you in your "column"?**
 - When they respond with a value N , respond $(N + 1)$ to the person who asked me.



How many students are in the lecture hall?

- We'll focus on solving the problem for single "column" of students.
 - I go to the first person in the front row and ask: "How many people are sitting directly behind you in your 'column'?"
 - Student's algorithm:
 - If there is no one behind me, answer 0.
 - **If someone is sitting behind me:**
 - **Ask that person: How many people are sitting directly behind you in your "column"?**
 - When they respond with a value N , respond $(N + 1)$ to the person who asked me.



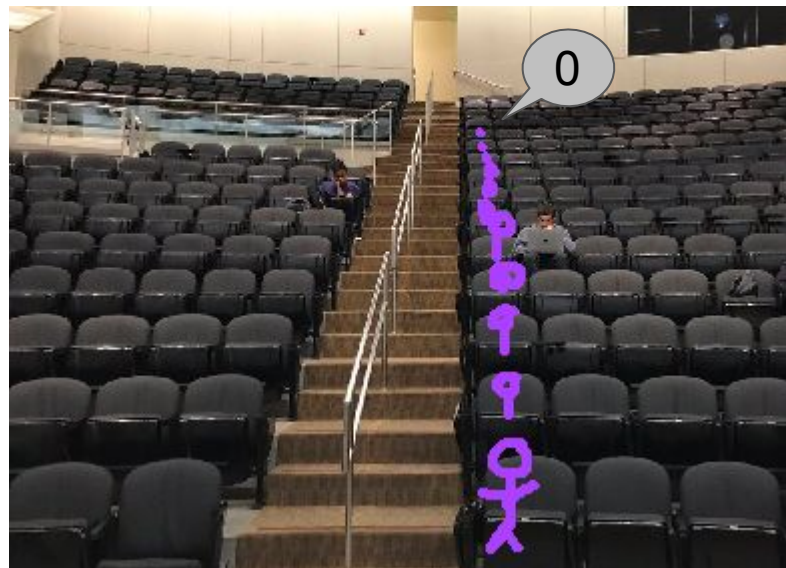
How many students are in the lecture hall?

- We'll focus on solving the problem for single "column" of students.
 - I go to the first person in the front row and ask: "How many people are sitting directly behind you in your 'column'?"
 - Student's algorithm:
 - If there is no one behind me, answer 0.
 - **If someone is sitting behind me:**
 - **Ask that person: How many people are sitting directly behind you in your "column"?**
 - When they respond with a value N , respond $(N + 1)$ to the person who asked me.



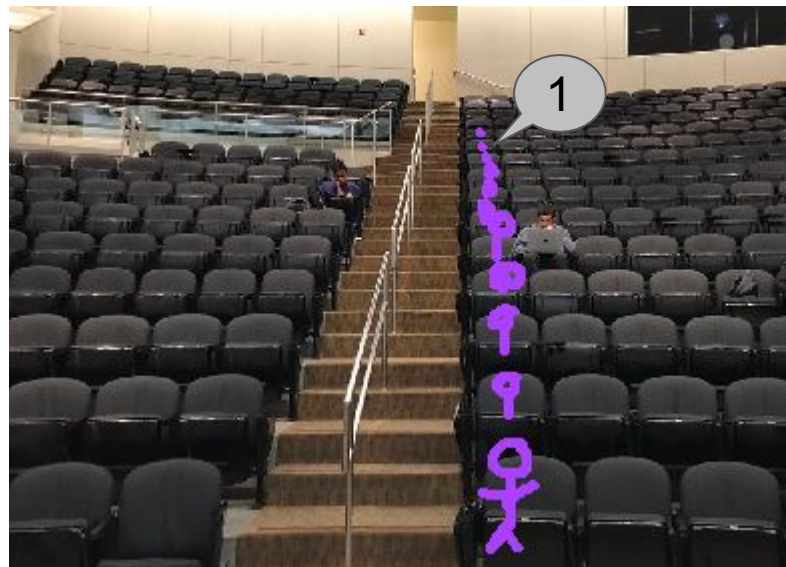
How many students are in the lecture hall?

- We'll focus on solving the problem for single "column" of students.
 - I go to the first person in the front row and ask: "How many people are sitting directly behind you in your 'column'?"
 - Student's algorithm:
 - **If there is no one behind me, answer 0.**
 - If someone is sitting behind me:
 - Ask that person: How many people are sitting directly behind you in your "column"?
 - When they respond with a value N , respond $(N + 1)$ to the person who asked me.



How many students are in the lecture hall?

- We'll focus on solving the problem for single "column" of students.
 - I go to the first person in the front row and ask: "How many people are sitting directly behind you in your 'column'?"
 - Student's algorithm:
 - If there is no one behind me, answer 0.
 - **If someone is sitting behind me:**
 - Ask that person: How many people are sitting directly behind you in your "column"?
 - When they respond with a value N , respond $(N + 1)$ to the person who asked me.



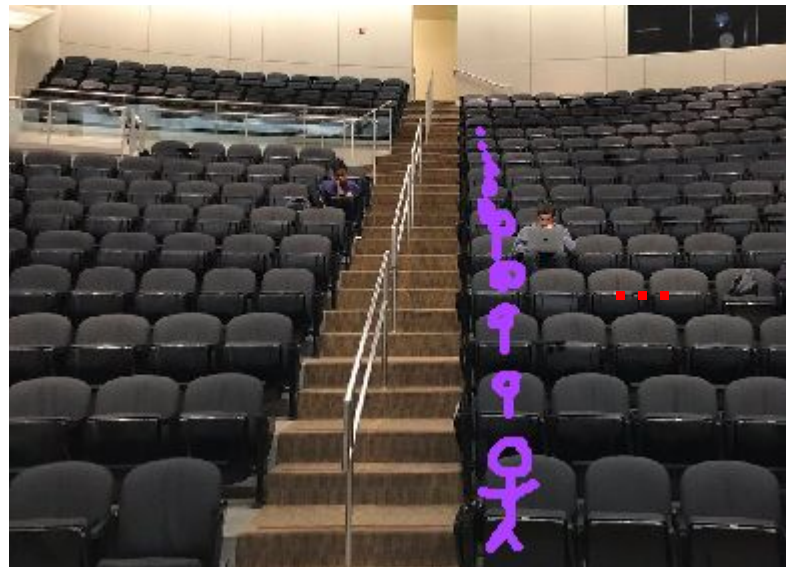
How many students are in the lecture hall?

- We'll focus on solving the problem for single "column" of students.
 - I go to the first person in the front row and ask: "How many people are sitting directly behind you in your 'column'?"
 - Student's algorithm:
 - If there is no one behind me, answer 0.
 - **If someone is sitting behind me:**
 - Ask that person: How many people are sitting directly behind you in your "column"?
 - When they respond with a value N , respond $(N + 1)$ to the person who asked me.



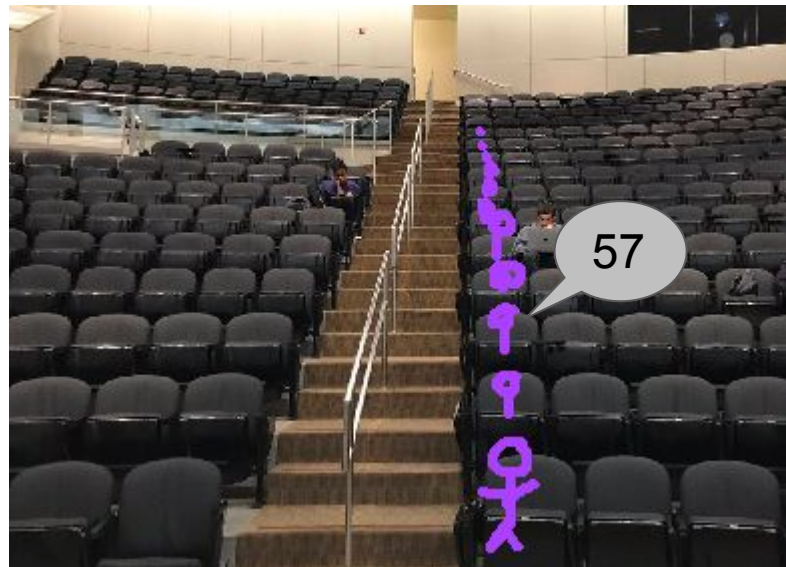
How many students are in the lecture hall?

- We'll focus on solving the problem for single “column” of students.
 - I go to the first person in the front row and ask: “How many people are sitting directly behind you in your ‘column’?”
 - Student's algorithm:
 - If there is no one behind me, answer 0.
 - **If someone is sitting behind me:**
 - Ask that person: How many people are sitting directly behind you in your “column”?
 - When they respond with a value N , respond $(N + 1)$ to the person who asked me.



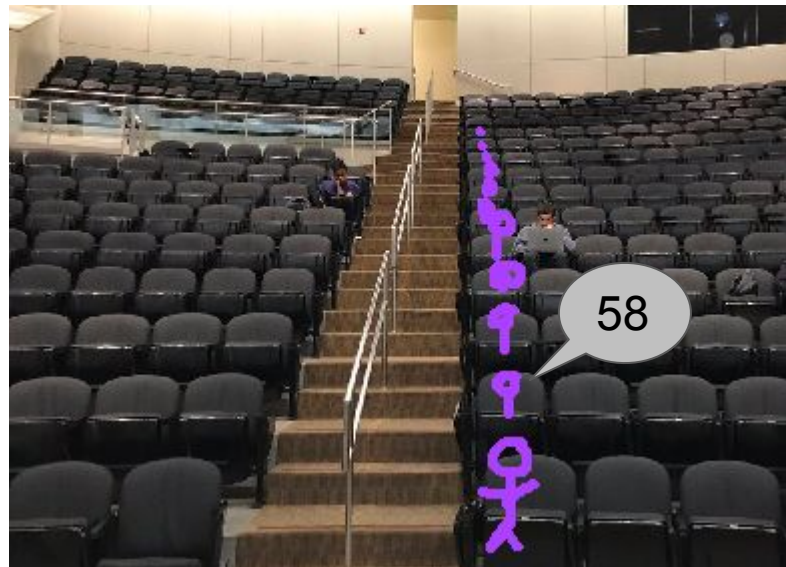
How many students are in the lecture hall?

- We'll focus on solving the problem for single "column" of students.
 - I go to the first person in the front row and ask: "How many people are sitting directly behind you in your 'column'?"
 - Student's algorithm:
 - If there is no one behind me, answer 0.
 - **If someone is sitting behind me:**
 - Ask that person: How many people are sitting directly behind you in your "column"?
 - When they respond with a value N , respond $(N + 1)$ to the person who asked me.



How many students are in the lecture hall?

- We'll focus on solving the problem for single "column" of students.
 - I go to the first person in the front row and ask: "How many people are sitting directly behind you in your 'column'?"
 - Student's algorithm:
 - If there is no one behind me, answer 0.
 - **If someone is sitting behind me:**
 - Ask that person: How many people are sitting directly behind you in your "column"?
 - When they respond with a value N , respond $(N + 1)$ to the person who asked me.



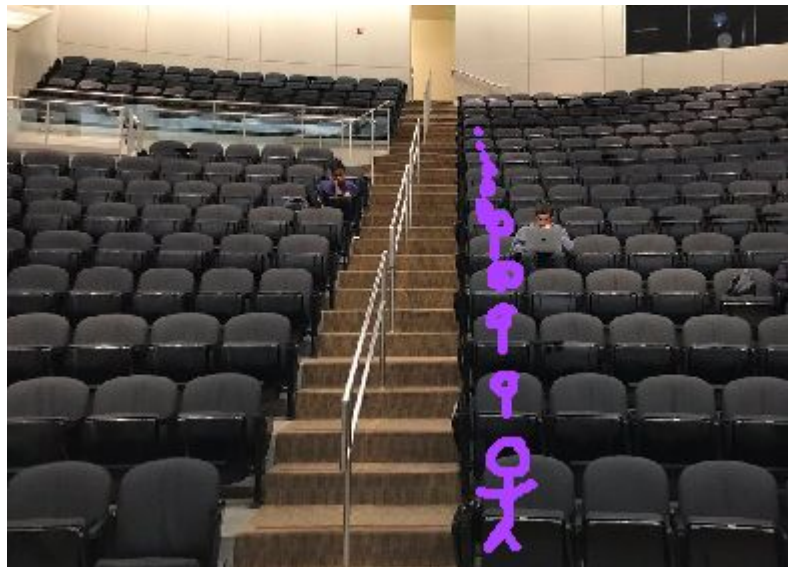
How many students are in the lecture hall?

- We'll focus on solving the problem for single "column" of students.
 - I go to the first person in the front row and ask: "How many people are sitting directly behind you in your 'column'?"
 - Student's algorithm:
 - If there is no one behind me, answer 0.
 - **If someone is sitting behind me:**
 - Ask that person: How many people are sitting directly behind you in your "column"?
 - When they respond with a value N , respond $(N + 1)$ to the person who asked me.



How many students are in the lecture hall?

- We'll focus on solving the problem for single "column" of students.
 - I go to the first person in the front row and ask: "How many people are sitting directly behind you in your 'column'?"
 - Student's algorithm:
 - If there is no one behind me, answer 0.
 - If someone is sitting behind me:
 - Ask that person: How many people are sitting directly behind you in your "column"?
 - When they respond with a value N , respond $(N + 1)$ to the person who asked me.
- Can generalize to the entire lecture hall!



Definition

recursion

A problem-solving technique in which tasks are completed by reducing them into repeated, smaller tasks of the same form.

Two main cases (components) of recursion

- Base case
 - The simplest version(s) of your problem that all other cases reduce to
 - An occurrence that can be answered directly

Two main cases (components) of recursion

- Base case
 - The simplest version(s) of your problem that all other cases reduce to
 - An occurrence that can be answered directly

“If there is no one behind me, answer 0.”

Two main cases (components) of recursion

- Base case
 - The simplest version(s) of your problem that all other cases reduce to
 - An occurrence that can be answered directly
- Recursive case
 - The step at which you break down more complex versions of the task into smaller occurrences
 - Cannot be answered directly
 - Take the “recursive leap of faith” and trust the smaller tasks will solve the problem for you!

Two main cases (components) of recursion

- Base case
 - The simplest version(s) of your problem that all other cases reduce to
 - An occurrence that can be answered directly
- Recursive case
 - The step at which you break down more complex versions of the task into smaller occurrences
 - Cannot be answered directly
 - Take the “recursive leap of faith” and trust the smaller tasks will solve the problem for you!

“If someone is sitting behind me...”

Two main cases (components) of recursion

- Base case
 - The simplest version(s) of your problem that all other cases reduce to
 - An occurrence that can be answered directly
- Recursive case
 - The step at which you break down more complex versions of the task into smaller occurrences
 - Cannot be answered directly
 - Take the “recursive leap of faith” and trust the smaller tasks will solve the problem for you!

Announcements

Announcements

- Assignment 2 is due Wednesday, 7/8.
- Assignment 3 will be released by the end of the day on Thursday.
- The mid-quarter diagnostic will cover through the end of this week (Thursday will be the last day of content covered).
- Please remember to only ask questions in the chat that are necessary for your immediate understanding!

Factorial example

Factorials

- The number **n factorial**, denoted **n!**, is

$$n \times (n - 1) \times \dots \times 3 \times 2 \times 1$$

Factorials

- The number **n factorial**, denoted **n!**, is

$$n \times (n - 1) \times \dots \times 3 \times 2 \times 1$$

- For example,
 - $3! = 3 \times 2 \times 1 = 6.$
 - $4! = 4 \times 3 \times 2 \times 1 = 24.$
 - $5! = 5 \times 4 \times 3 \times 2 \times 1 = 120.$
 - $0! = 1.$ (by definition)

Factorials

- The number **n factorial**, denoted **n!**, is

$$n \times (n - 1) \times \dots \times 3 \times 2 \times 1$$

- For example,
 - $3! = 3 \times 2 \times 1 = 6.$
 - $4! = 4 \times 3 \times 2 \times 1 = 24.$
 - $5! = 5 \times 4 \times 3 \times 2 \times 1 = 120.$
 - $0! = 1.$ (by definition)
- Factorials show up in unexpected places. We'll see one later this quarter when we talk about sorting algorithms.

Factorials

- The number **n factorial**, denoted **n!**, is

$$n \times (n - 1) \times \dots \times 3 \times 2 \times 1$$

- For example,
 - $3! = 3 \times 2 \times 1 = 6.$
 - $4! = 4 \times 3 \times 2 \times 1 = 24.$
 - $5! = 5 \times 4 \times 3 \times 2 \times 1 = 120.$
 - $0! = 1.$ (by definition)
- Factorials show up in unexpected places. We'll see one later this quarter when we talk about sorting algorithms.
- Let's implement a function to compute factorials!

Computing factorials

$$5! = 5 \times 4 \times 3 \times 2 \times 1$$

Computing factorials

$$5! = 5 \times 4 \times 3 \times 2 \times 1$$

Computing factorials

$$5! = 5 \times \underbrace{4 \times 3 \times 2 \times 1}_{4!}$$

Computing factorials

$$5! = 5 \times 4!$$

Computing factorials

$$5! = 5 \times 4!$$

Computing factorials

$$5! = 5 \times 4!$$

$$4! = 4 \times 3 \times 2 \times 1$$

Computing factorials

$$5! = 5 \times 4!$$

$$4! = 4 \times 3 \times 2 \times 1$$

Computing factorials

$$5! = 5 \times 4!$$

$$4! = 4 \times \underbrace{3 \times 2 \times 1}_{3!}$$

Computing factorials

$$5! = 5 \times 4!$$

$$4! = 4 \times 3!$$

Computing factorials

$$5! = 5 \times 4!$$

$$4! = 4 \times 3!$$

Computing factorials

$$5! = 5 \times 4!$$

$$4! = 4 \times 3!$$

$$3! = 3 \times 2 \times 1$$

Computing factorials

$$5! = 5 \times 4!$$

$$4! = 4 \times 3!$$

$$3! = 3 \times 2 \times 1$$

Computing factorials

$$5! = 5 \times 4!$$

$$4! = 4 \times 3!$$

$$3! = 3 \times \underbrace{2 \times 1}_{2!}$$

Computing factorials

$$5! = 5 \times 4!$$

$$4! = 4 \times 3!$$

$$3! = 3 \times 2!$$

Computing factorials

$$5! = 5 \times 4!$$

$$4! = 4 \times 3!$$

$$3! = 3 \times 2!$$

Computing factorials

$$5! = 5 \times 4!$$

$$4! = 4 \times 3!$$

$$3! = 3 \times 2!$$

$$2! = 2 \times 1!$$

Computing factorials

$$5! = 5 \times 4!$$

$$4! = 4 \times 3!$$

$$3! = 3 \times 2!$$

$$2! = 2 \times 1!$$

$$1! = 1 \times 0!$$

Computing factorials

$$5! = 5 \times 4!$$

$$4! = 4 \times 3!$$

$$3! = 3 \times 2!$$

$$2! = 2 \times 1!$$

$$1! = 1 \times 0!$$

$$0! = 1$$

Computing factorials

$$5! = 5 \times 4!$$

$$4! = 4 \times 3!$$

$$3! = 3 \times 2!$$

$$2! = 2 \times 1!$$

$$1! = 1 \times 0!$$

$$0! = 1$$

By definition!



Another view of factorials

$$n! = \begin{cases} 1 & \text{if } n = 0 \\ n \times (n - 1)! & \text{otherwise} \end{cases}$$

Another view of factorials

$$n! = \begin{cases} 1 & \text{if } n = 0 \\ n \times (n - 1)! & \text{otherwise} \end{cases}$$

```
int factorial (int n) {  
    if (n == 0) {  
        return 1;  
    } else {  
        return n * factorial(n-1);  
    }  
}
```

Recursion in action

```
int main() {  
    int n = factorial(5);  
    cout << "5! = " << n << endl;  
    return 0;  
}
```

Recursion in action

```
int main() {  
    int n = factorial(5);  
    cout << "5! = " << n << endl;  
    return 0;  
}
```

This is a “**stack frame**.” One gets created each time a function is called.

- The “stack” is where in your computer’s memory the information is stored.
- A “frame” stores all of the data (variables) for that particular function call.

Recursion in action

```
int main() {  
    int n = factorial(5);  
    cout << "5! = " << n << endl;  
    return 0;  
}
```

Recursion in action

```
int main() {
```

```
    int factorial (int n) {  
        if (n == 0) {  
            return 1;  
        } else {  
            return n * factorial(n-1);  
        }  
    }  
}
```



n



When a function gets called, a new stack frame gets created.

Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
        if (n == 0) {
```

```
            return 1;
```

```
        } else {
```

```
            return n * factorial(n-1);
```

```
        }
```

```
    }
```



n

Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
        if (n == 0) {
```

```
            return 1;
```

```
        } else {
```

```
            return n * factorial(n-1);
```

```
        }
```

```
    }
```



n

Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
        if (n == 0) {
```

```
            return 1;
```

```
        } else {
```

```
            return n * factorial(n-1);
```

```
        }
```

```
    }
```



n

Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
        if (n == 0) {
```

```
            return 1;
```

```
        } else {
```

```
            return n * factorial(n-1);
```

```
        }
```

```
    }
```



n

5

Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
        if (n == 0) {
```

```
            return 1;
```

```
        } else {
```

```
            return n * factorial(n-1);
```

```
        }
```

```
    }
```

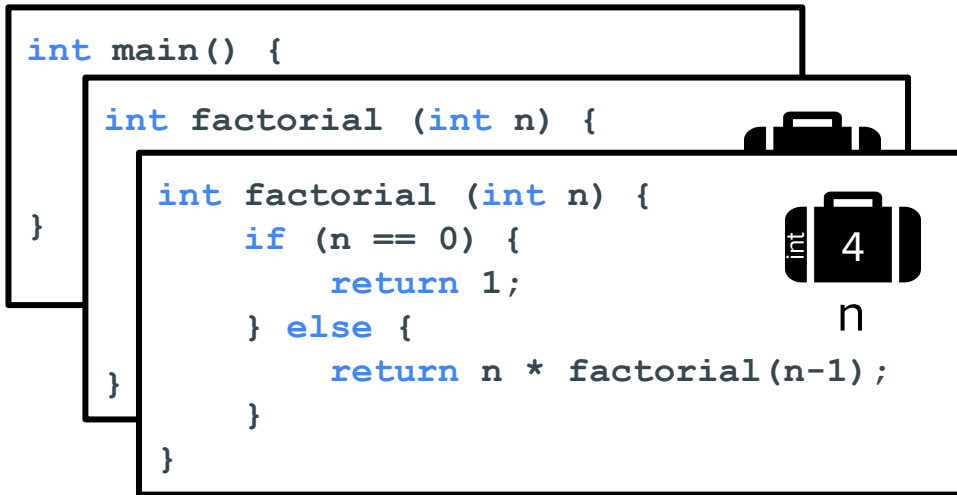


`n`

5

Recursion in action

```
int main() {  
    int factorial (int n) {  
        int factorial (int n) {  
            if (n == 0) {  
                return 1;  
            } else {  
                return n * factorial(n-1);  
            }  
        }  
    }  
}
```



Every time we call **factorial()**, we get a new copy of the local variable **n** that's independent of all the previous copies because it exists inside the new frame.



Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
    }
```

```
        int factorial (int n) {
```

```
            if (n == 0) {
```

```
                return 1;
```

```
            } else {
```

```
                return n * factorial(n-1);
```

```
            }
```

```
        }
```



n

Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
    }
```

```
        int factorial (int n) {
```

```
            if (n == 0) {
```

```
                return 1;
```

```
            } else {
```

```
                return n * factorial(n-1);
```

```
            }
```

```
        }
```



n

Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
        int factorial (int n) {
```

```
            if (n == 0) {
```

```
                return 1;
```

```
            } else {
```

```
                return n * factorial(n-1);
```

```
            }
```

```
        }
```



n

Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
    }
```

```
        int factorial (int n) {
```

```
            if (n == 0) {
```

```
                return 1;
```

```
            } else {
```

```
                return n * factorial(n-1);
```

```
            }
```

```
        }
```



n

Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
        int factorial (int n) {
```

```
            if (n == 0) {
```

```
                return 1;
```

```
            } else {
```

```
                return n * factorial(n-1);
```

```
            }
```

```
        }
```

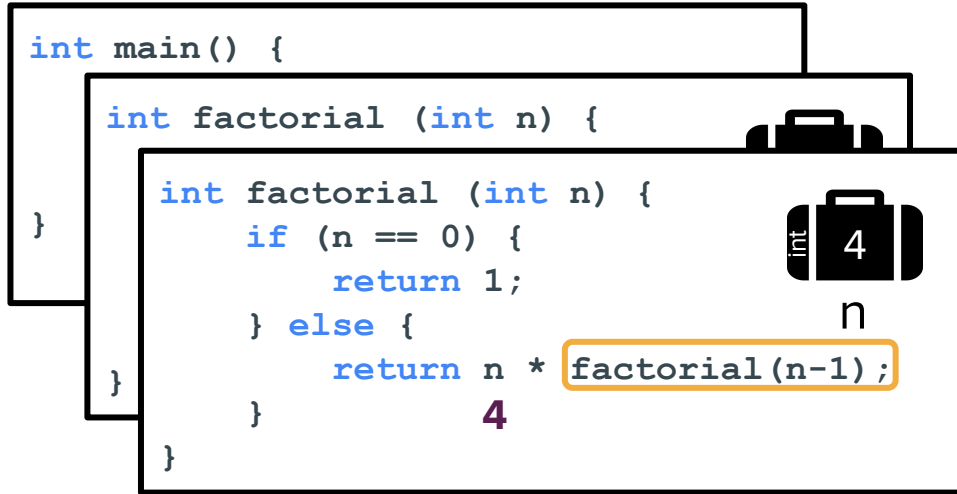


n

4

Recursion in action

```
int main() {  
    int factorial (int n) {  
        int factorial (int n) {  
            if (n == 0) {  
                return 1;  
            } else {  
                return n * factorial(n-1);  
            }  
        }  
    }  
}
```



Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
    }
```

```
        int factorial (int n) {
```

```
            int factorial (int n) {
```

```
                if (n == 0) {
```

```
                    return 1;
```

```
                } else {
```

```
                    return n * factorial(n-1);
```

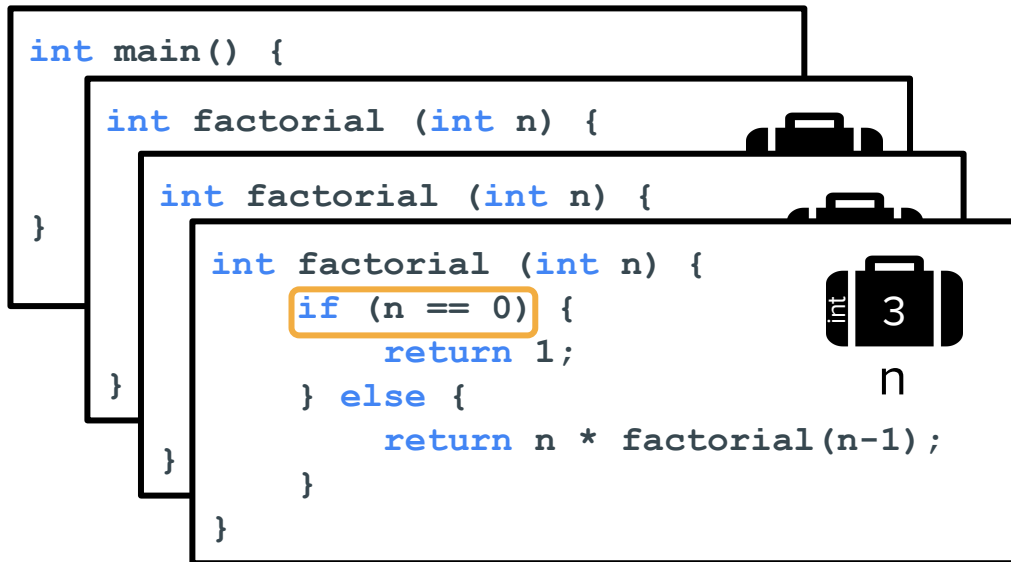
```
                }
```

```
            }
```

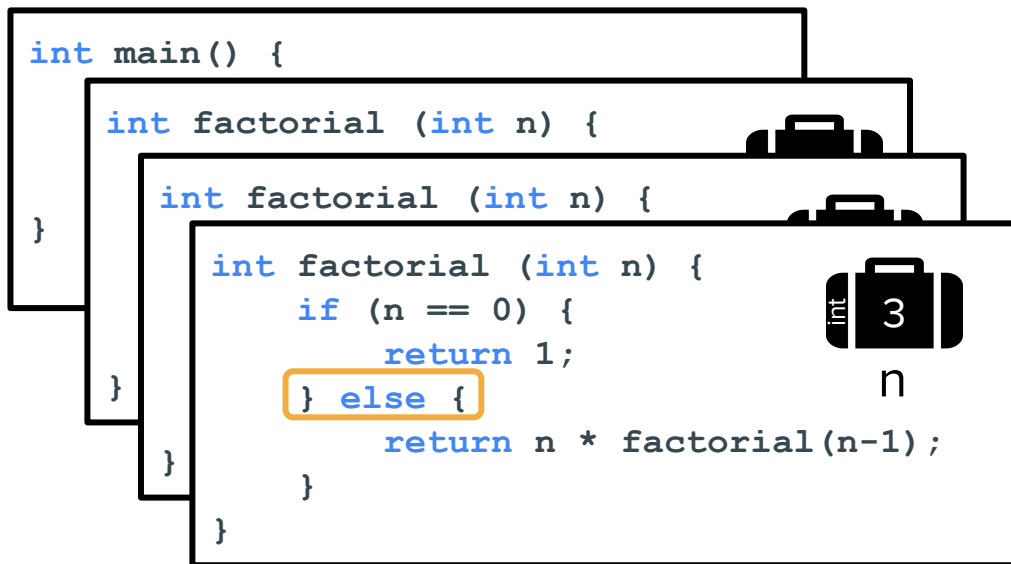


n

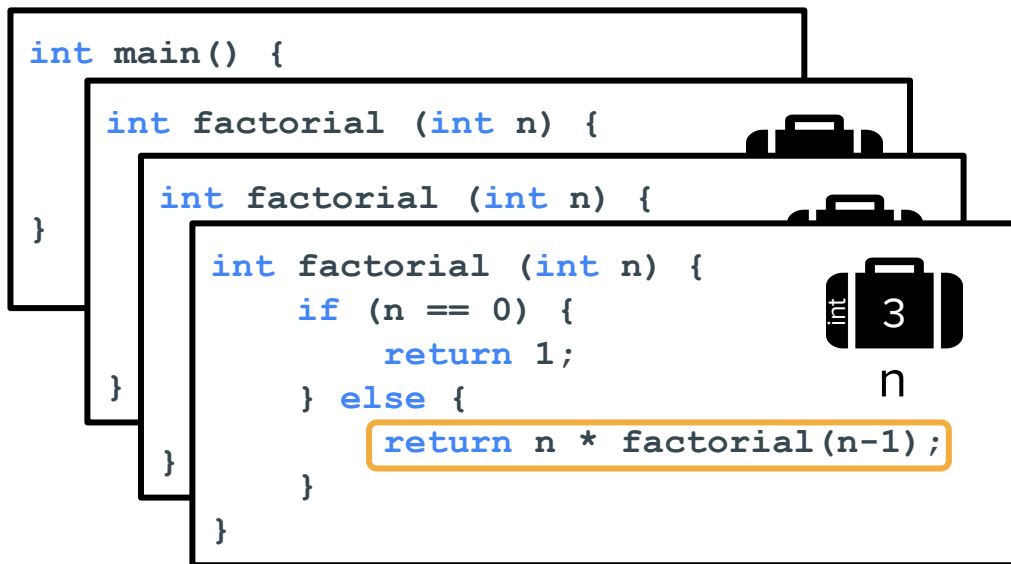
Recursion in action



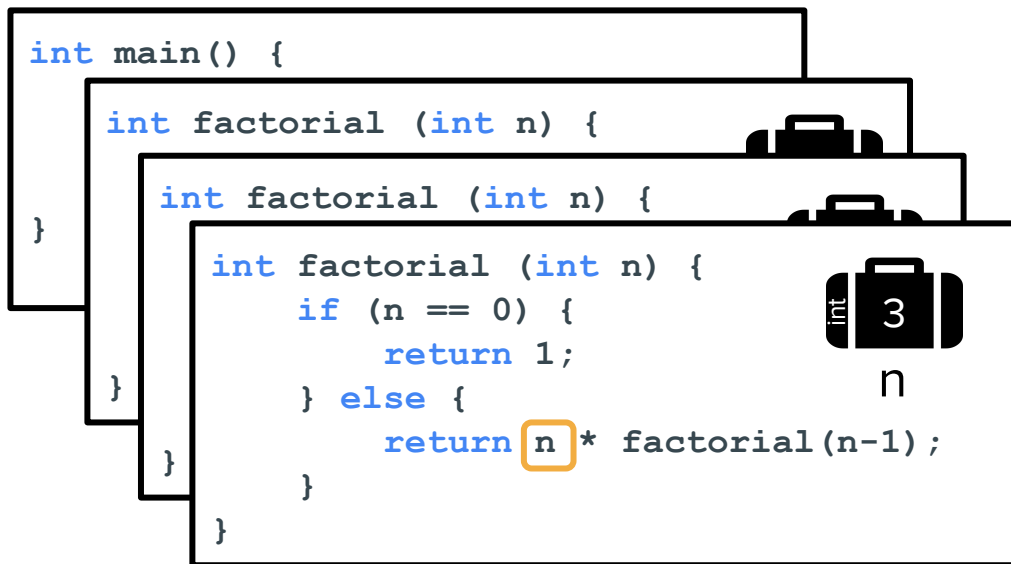
Recursion in action



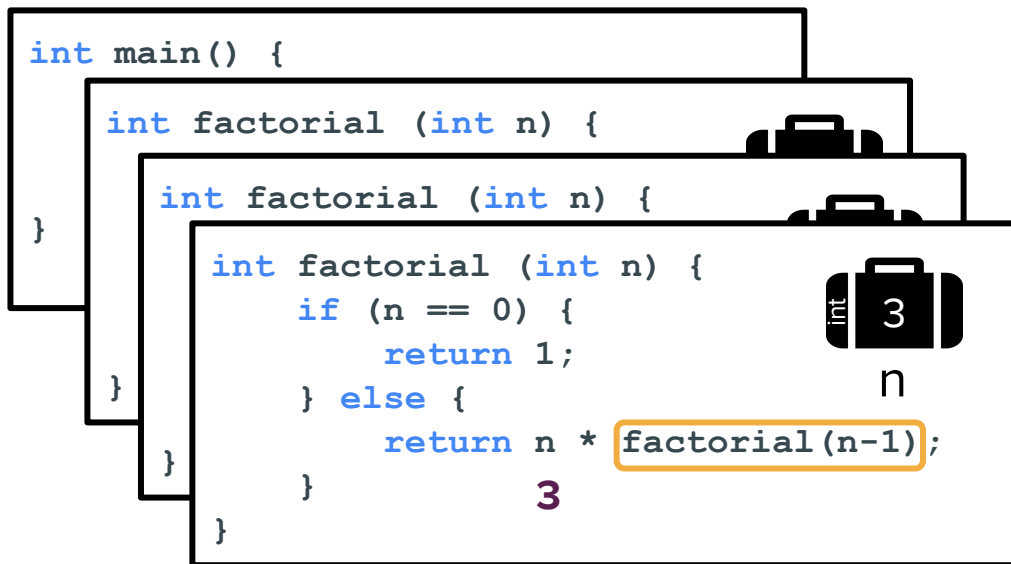
Recursion in action



Recursion in action



Recursion in action



Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
        int factorial (int n) {
```

```
            int factorial (int n) {
```

```
                int factorial (int n) {
```

```
                    if (n == 0) {
```

```
                        return 1;
```

```
                    } else {
```

```
                        return n * factorial(n-1);
```

```
                }
```

```
            }
```

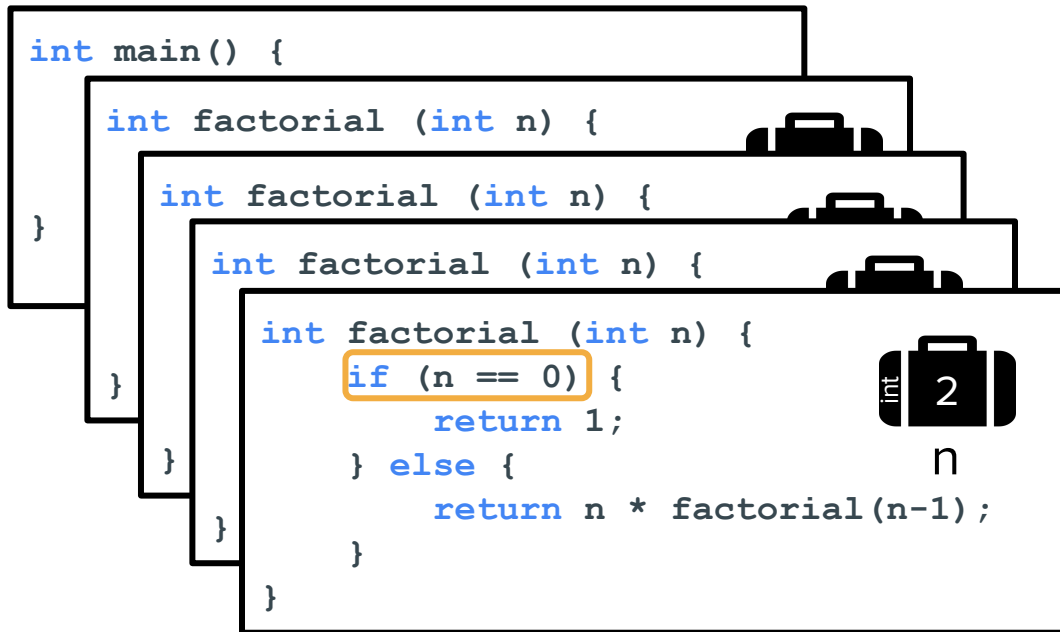
```
        }
```

```
    }
```

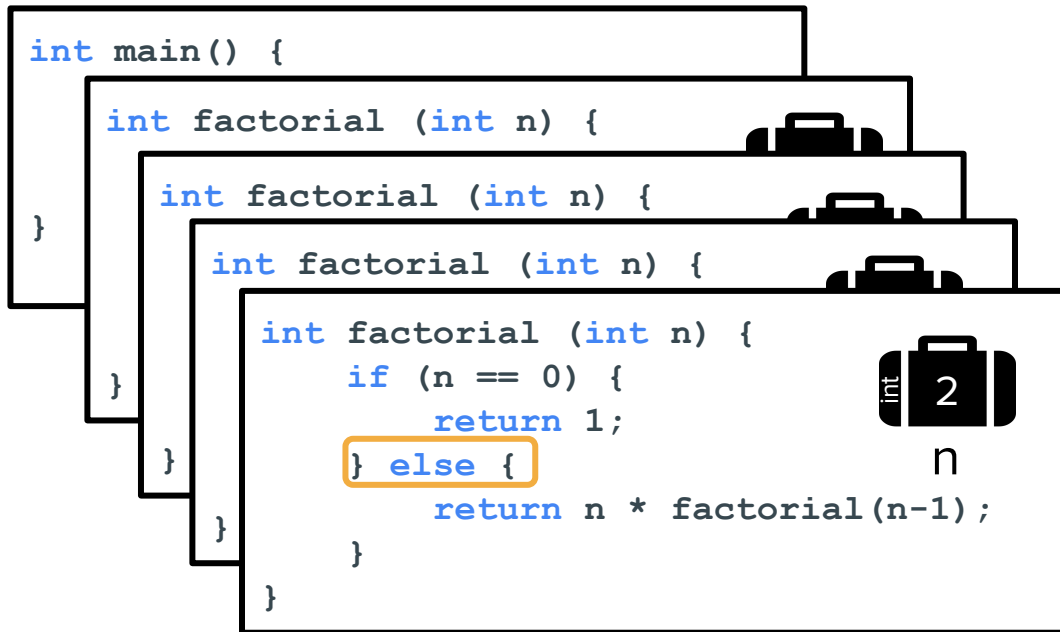


n

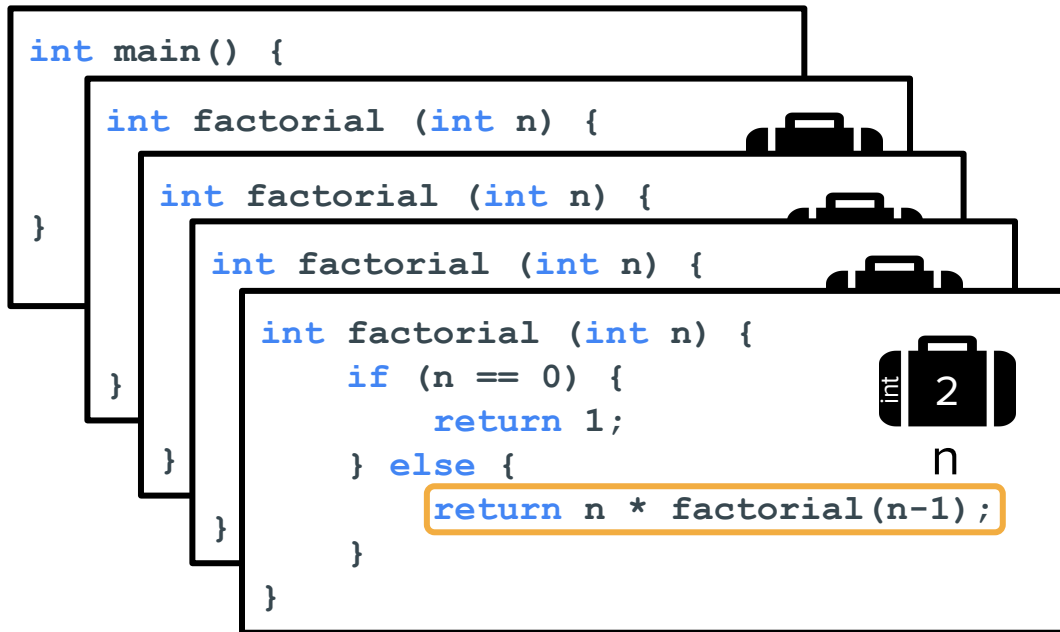
Recursion in action



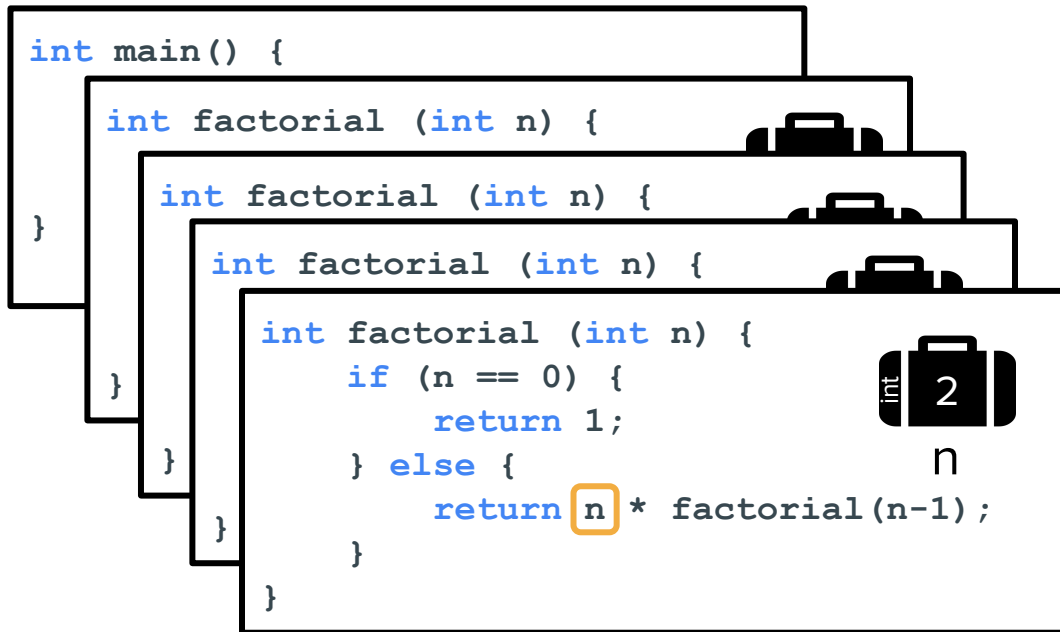
Recursion in action



Recursion in action



Recursion in action



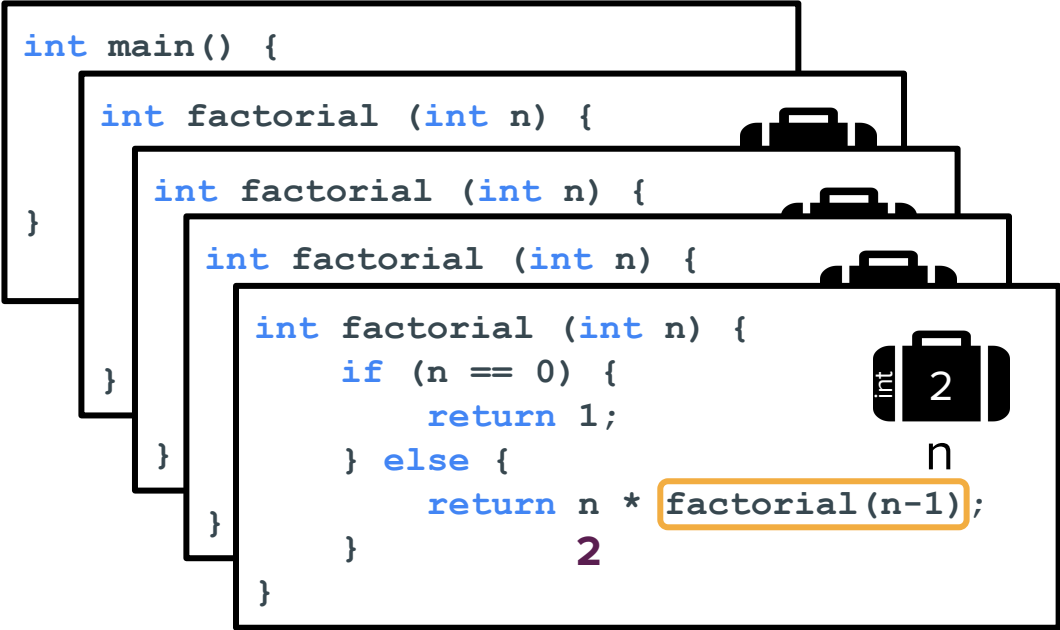
Recursion in action

```
int main() {  
  int factorial (int n) {  
    int factorial (int n) {  
      int factorial (int n) {  
        int factorial (int n) {  
          if (n == 0) {  
            return 1;  
          } else {  
            return n * factorial(n-1);  
          }  
        }  
      }  
    }  
  }  
}
```

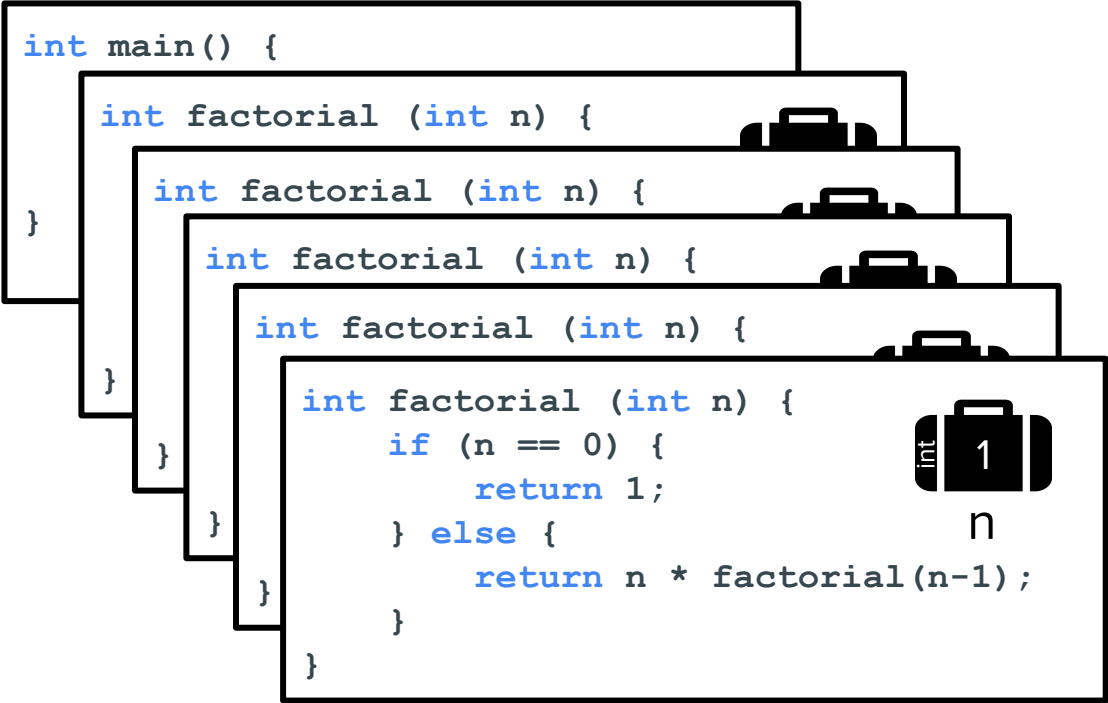
The diagram illustrates the recursive calls for calculating the factorial of 2. The code is shown in four overlapping boxes, representing the call stack. The innermost box shows the base case: `if (n == 0) { return 1; }`. The next box shows the recursive call for `n=1`: `} else { return 1 * factorial(1-1); }`. The next box shows the recursive call for `n=2`: `} else { return 2 * factorial(2-1); }`. The outermost box shows the `main` function: `int main() { int factorial (int n) { ... }`. A suitcase icon with 'int' and '2' is next to the return statement in the `n=2` box. The variable `n` is also labeled with `n`.

Recursion in action

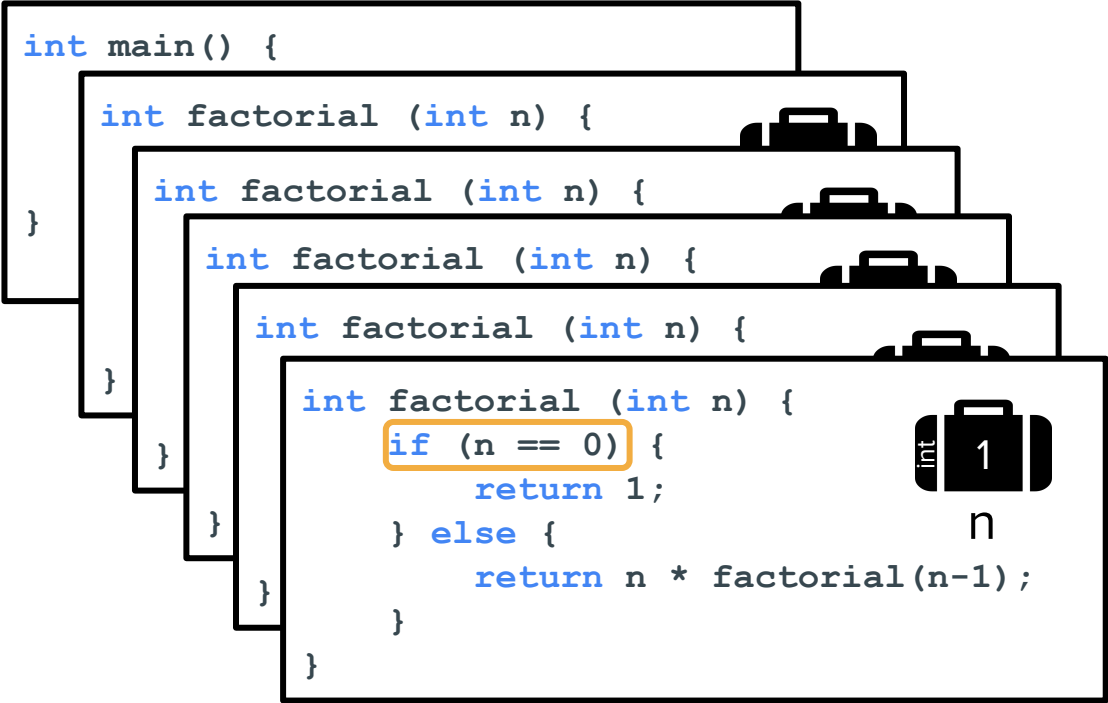
```
int main() {  
  int factorial (int n) {  
    int factorial (int n) {  
      int factorial (int n) {  
        int factorial (int n) {  
          if (n == 0) {  
            return 1;  
          } else {  
            return n * factorial(n-1);  
          }  
        }  
      }  
    }  
  }  
}
```



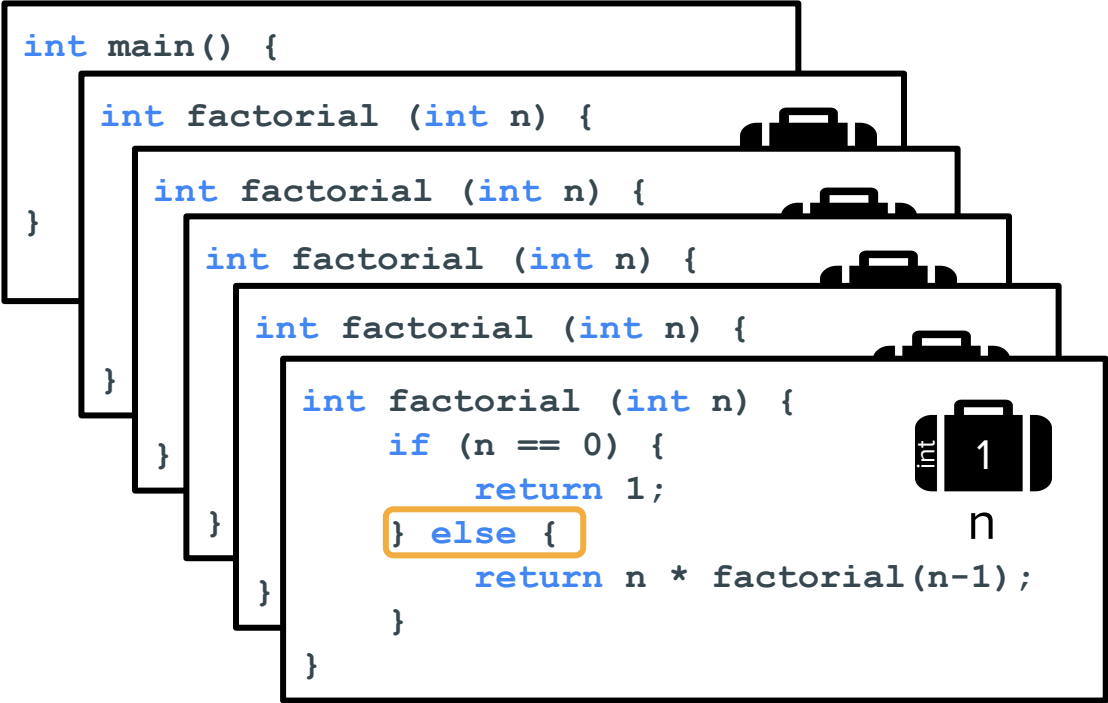
Recursion in action



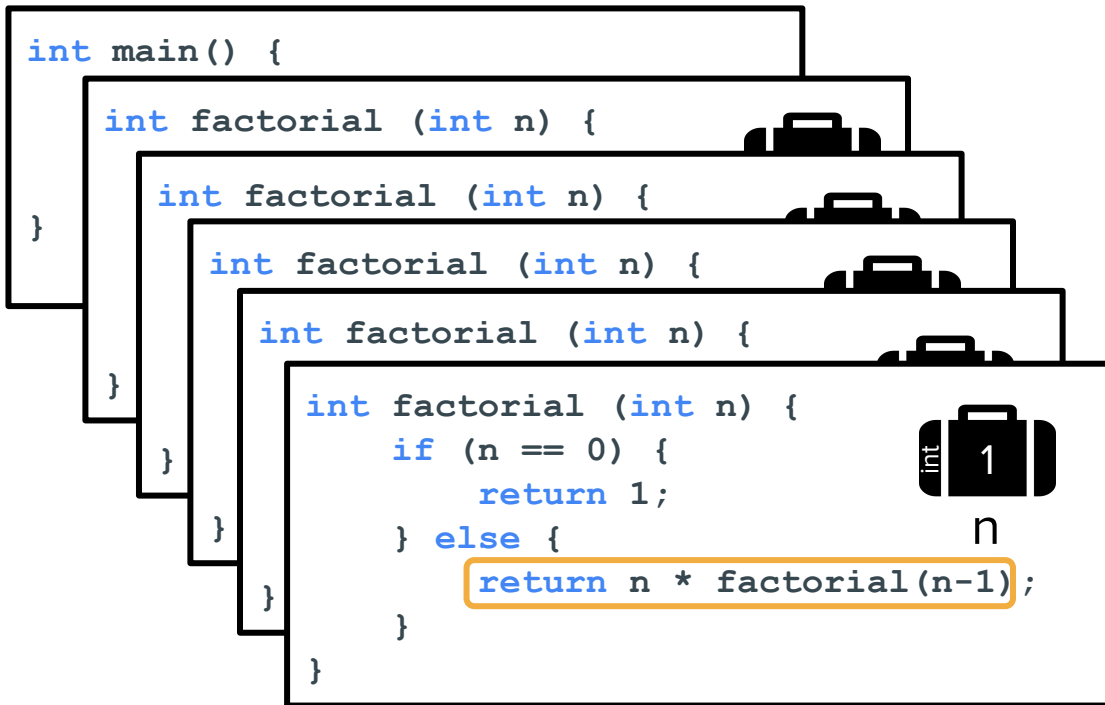
Recursion in action



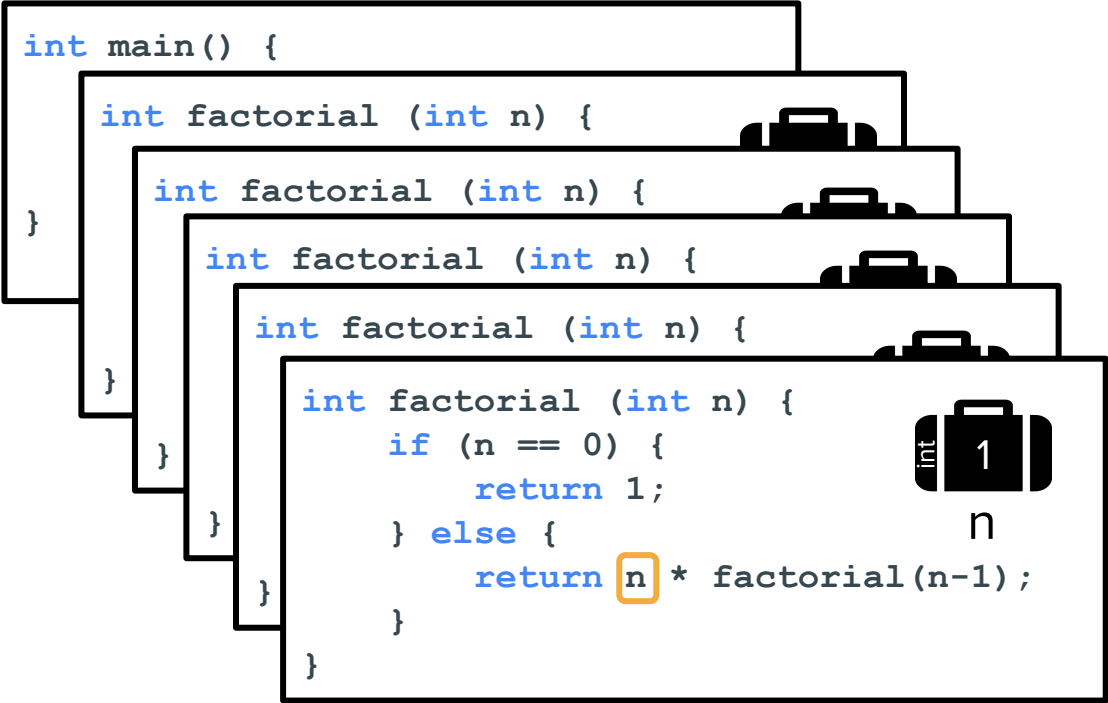
Recursion in action



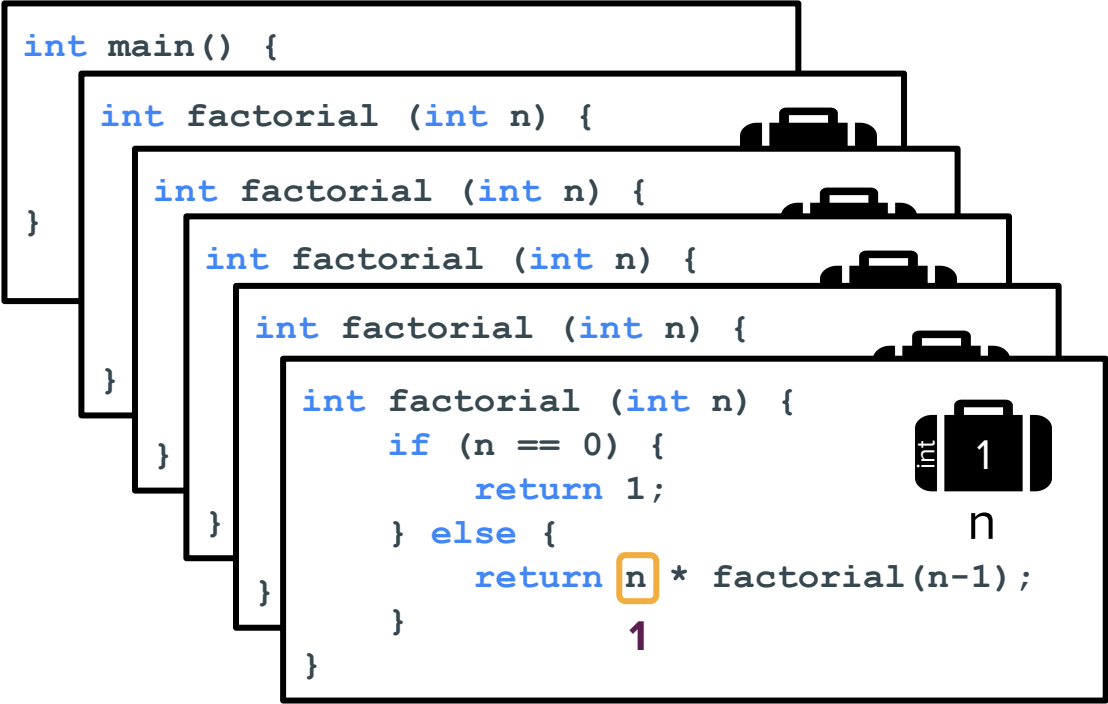
Recursion in action



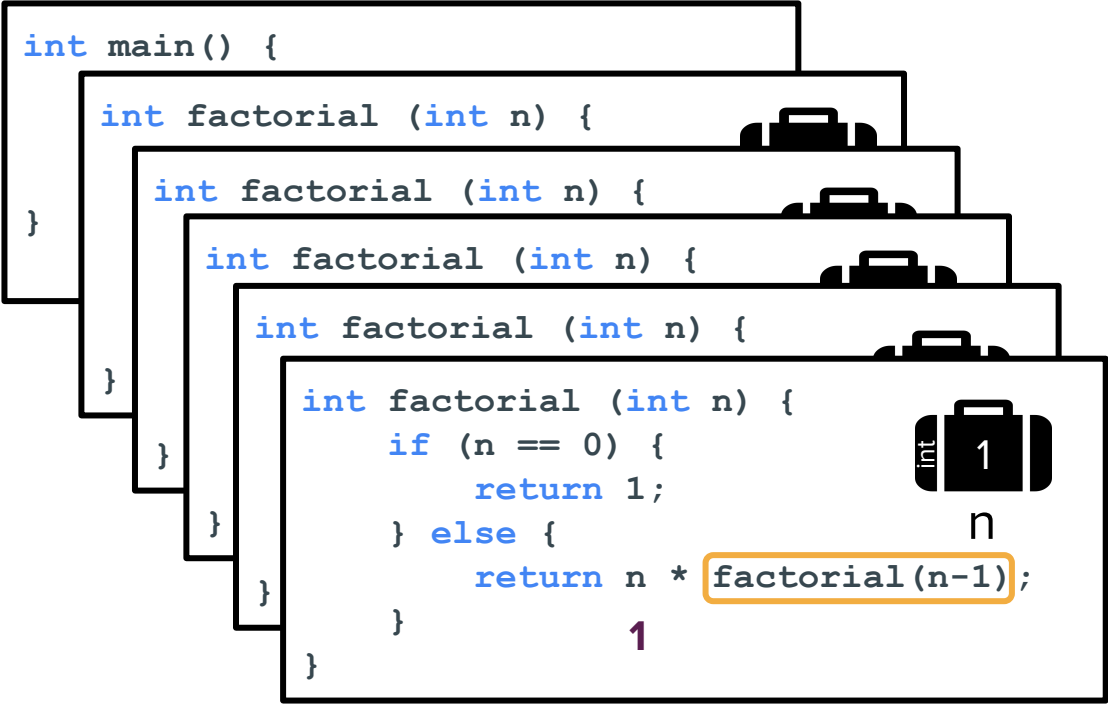
Recursion in action



Recursion in action



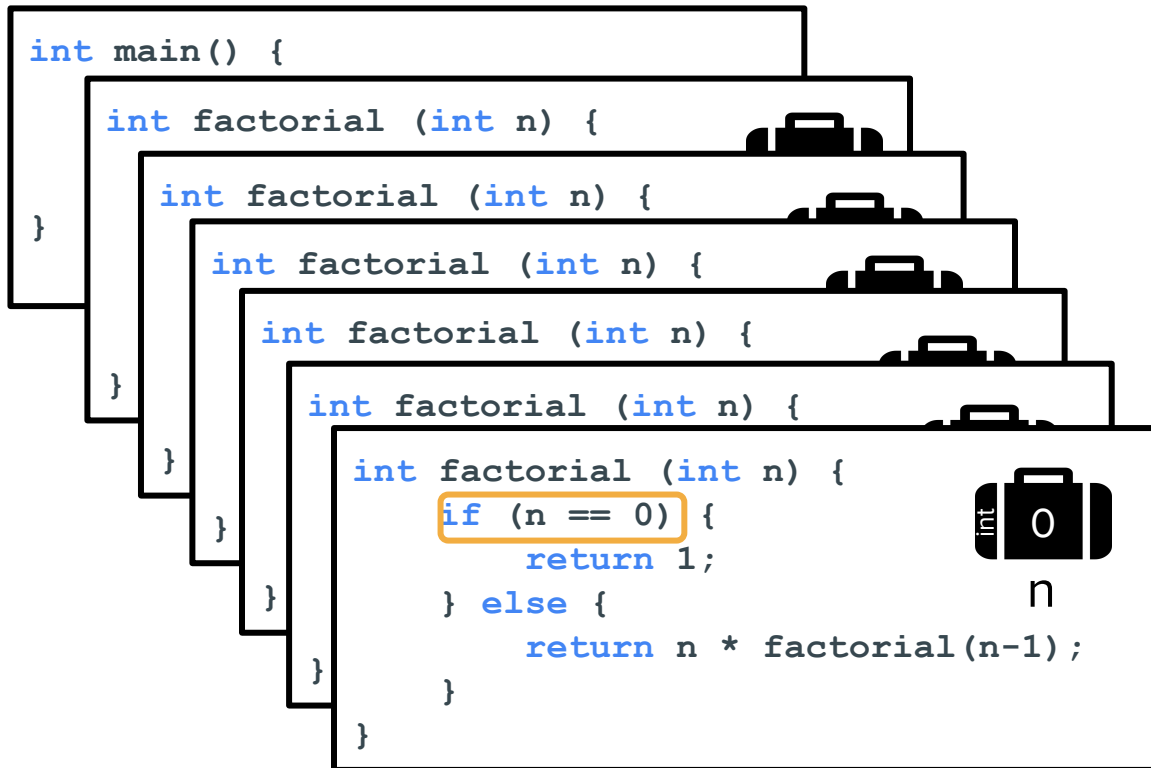
Recursion in action



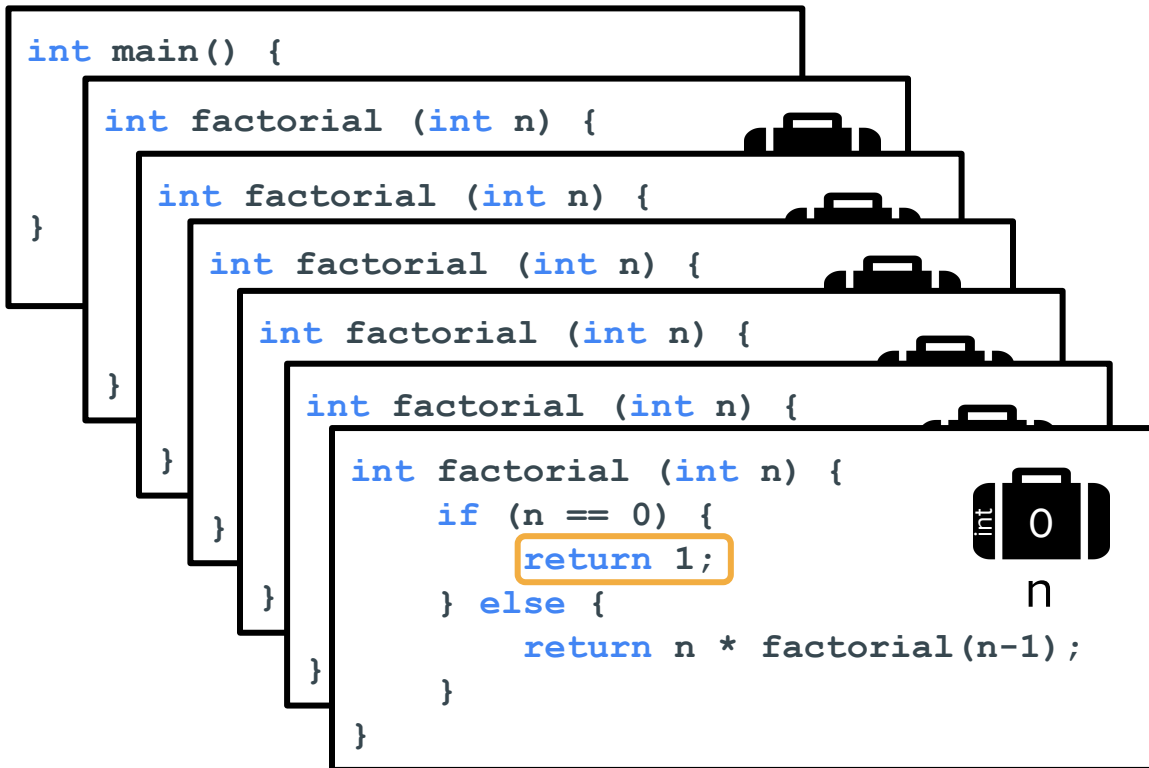
Recursion in action

```
int main() {  
  int factorial (int n) {  
    int factorial (int n) {  
      int factorial (int n) {  
        int factorial (int n) {  
          int factorial (int n) {  
            if (n == 0) {  
              return 1;  
            } else {  
              return n * factorial(n-1);  
            }  
          }  
        }  
      }  
    }  
  }  
}
```

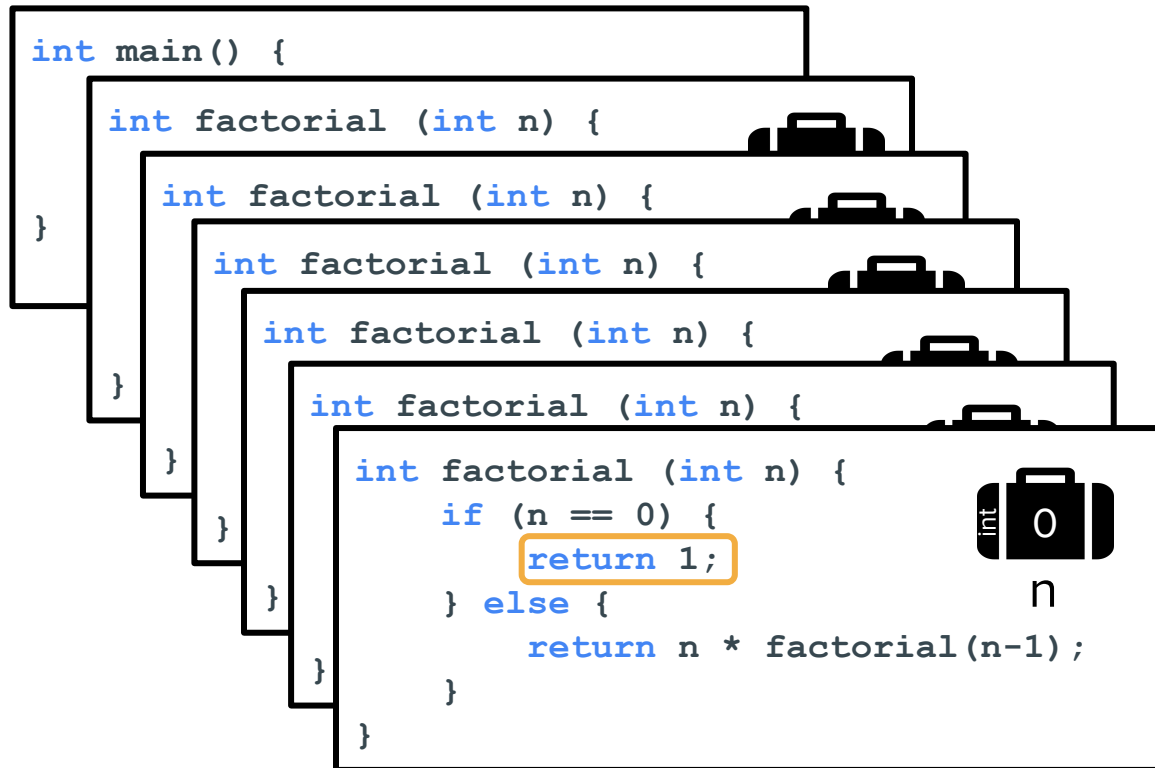
Recursion in action



Recursion in action



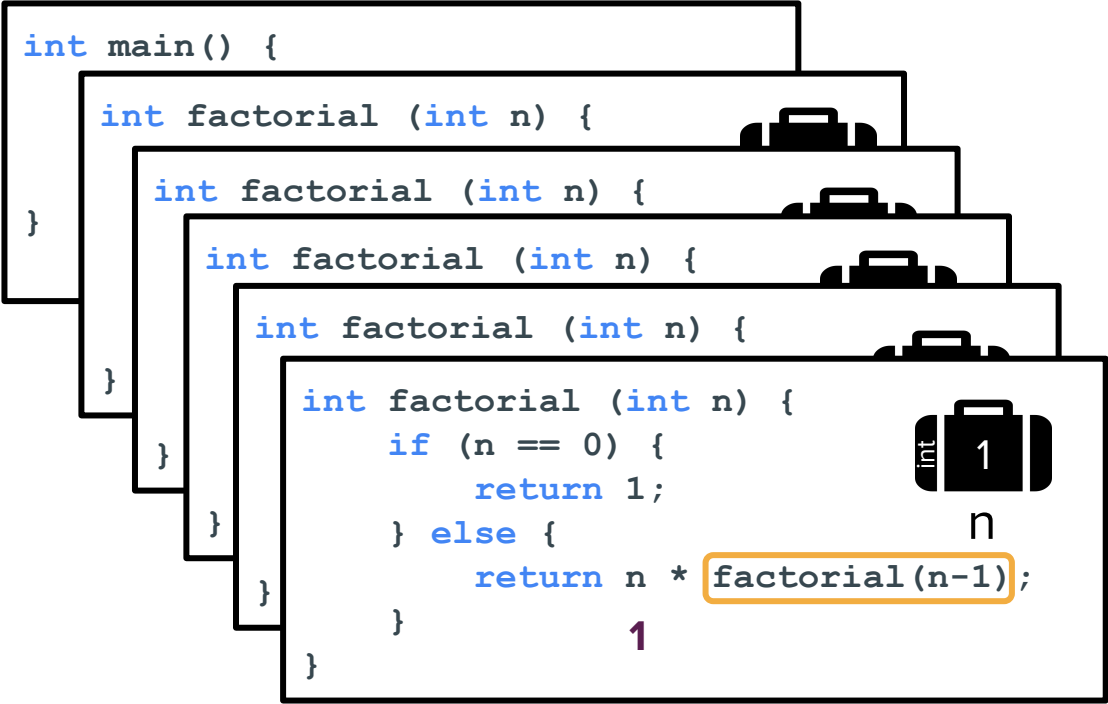
Recursion in action



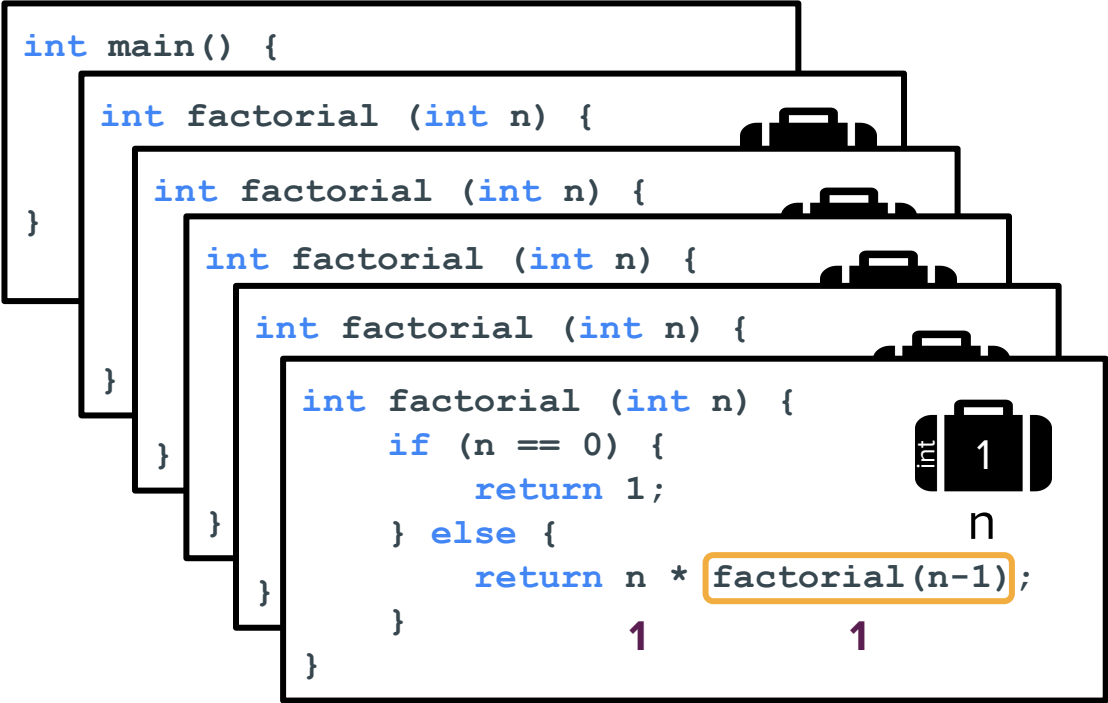
Stack frames go away (get cleared from memory) once they return.



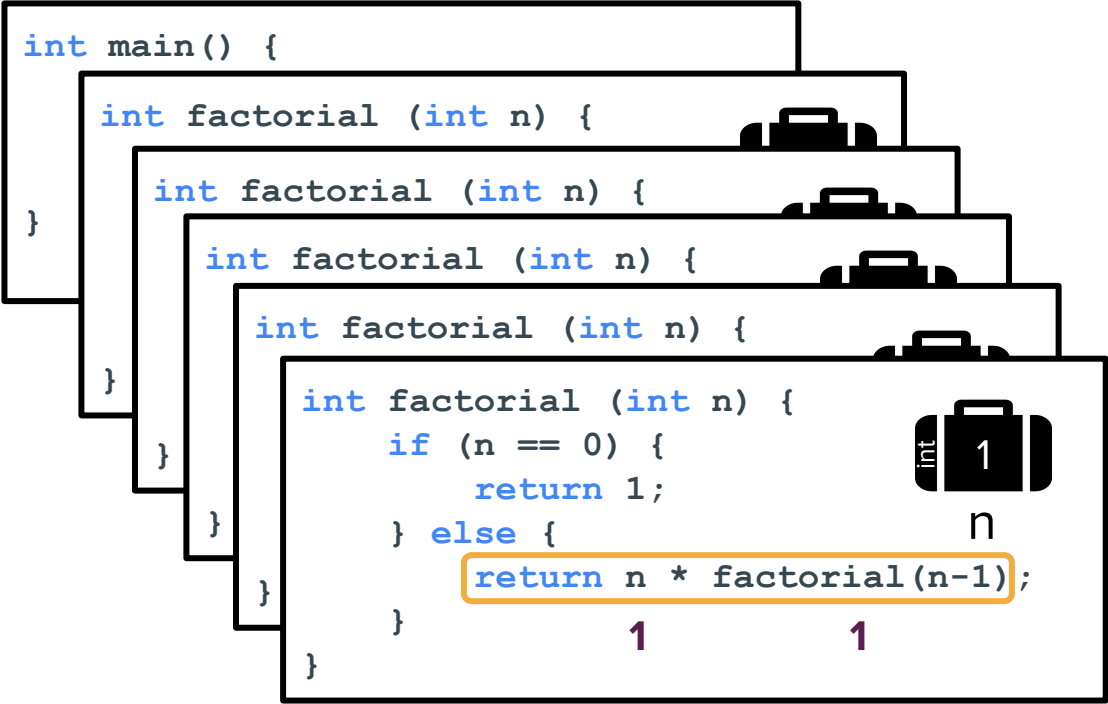
Recursion in action



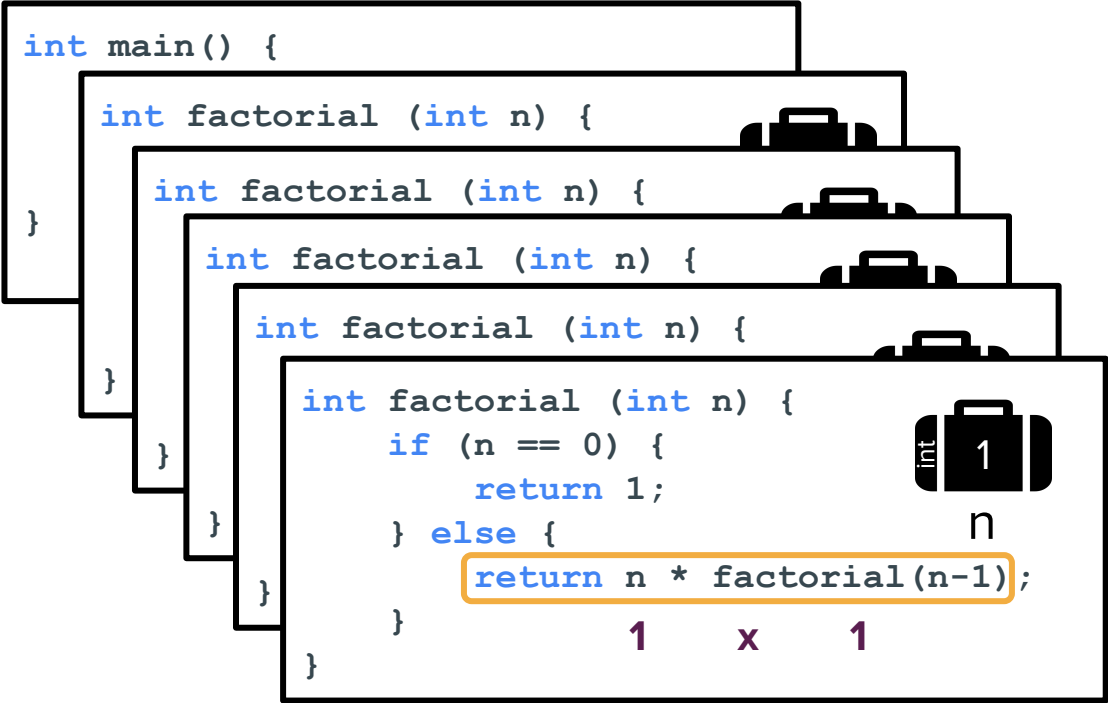
Recursion in action



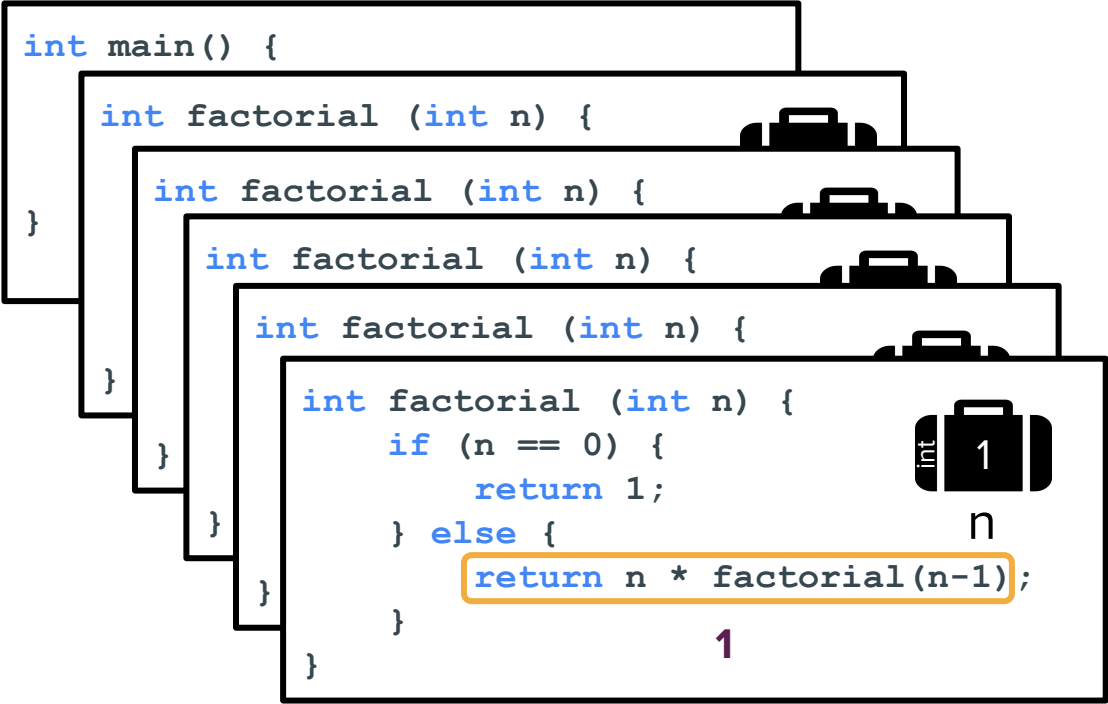
Recursion in action



Recursion in action



Recursion in action

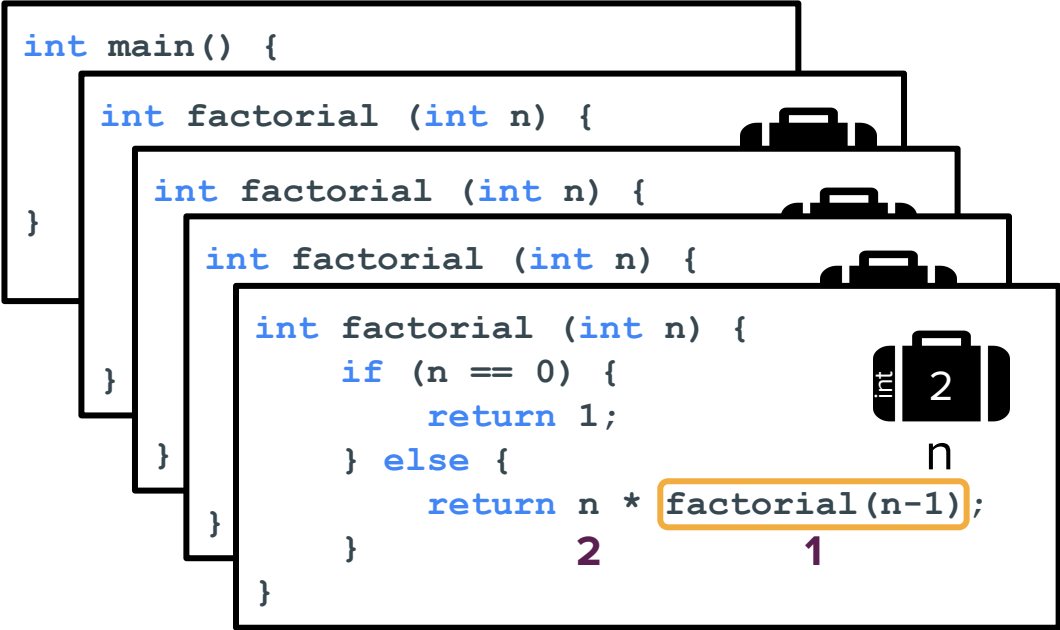


Recursion in action

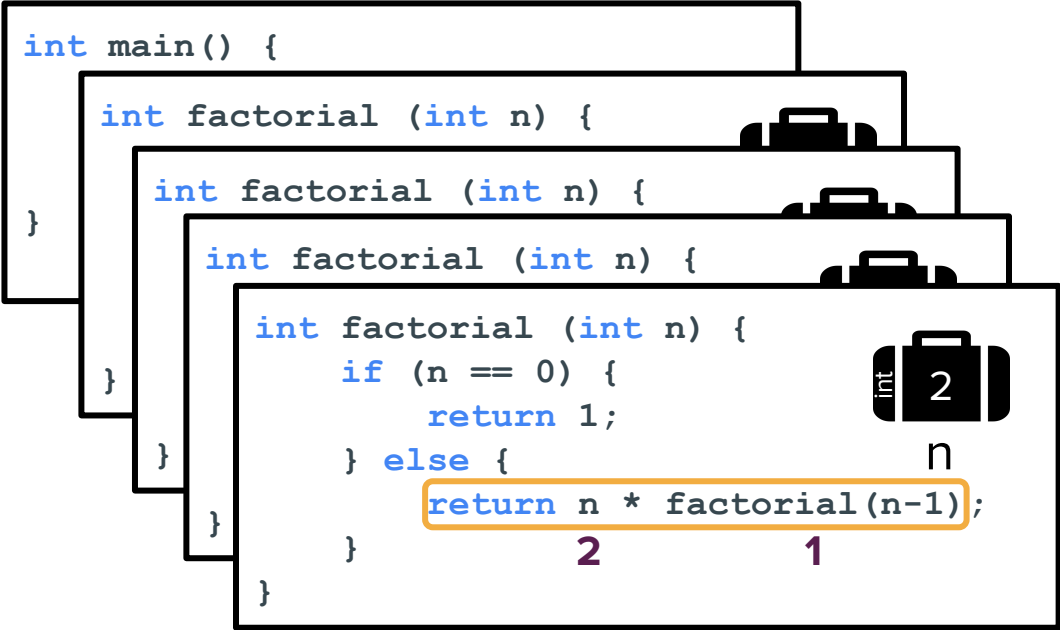
```
int main() {  
  int factorial (int n) {  
    int factorial (int n) {  
      int factorial (int n) {  
        int factorial (int n) {  
          if (n == 0) {  
            return 1;  
          } else {  
            return n * factorial(n-1);  
          }  
        }  
      }  
    }  
  }  
}
```

int 2
n

Recursion in action



Recursion in action



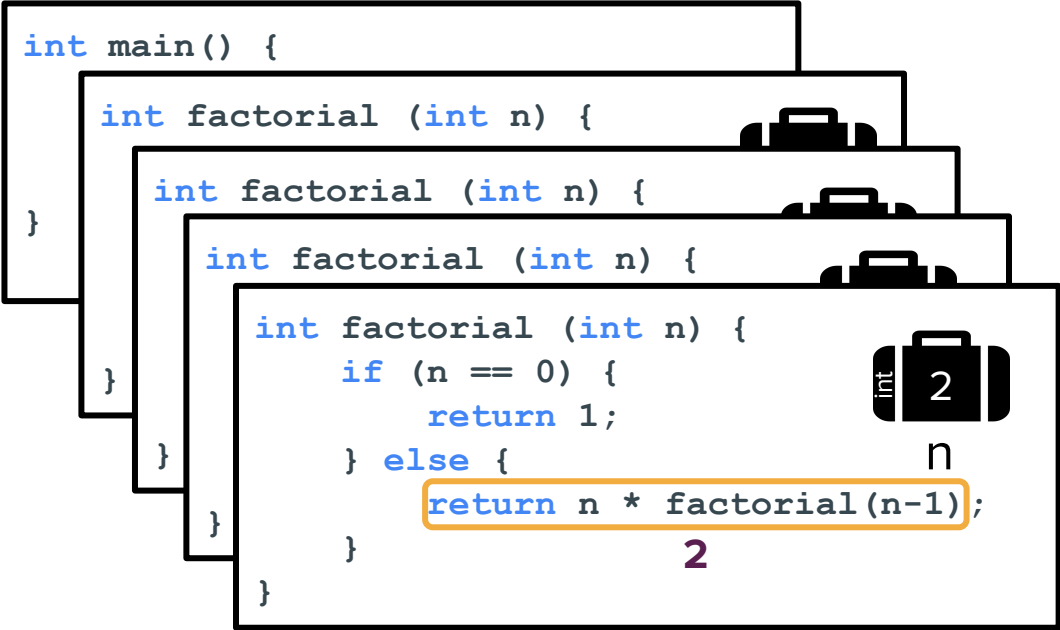
Recursion in action

```
int main() {  
  int factorial (int n) {  
    int factorial (int n) {  
      int factorial (int n) {  
        int factorial (int n) {  
          if (n == 0) {  
            return 1;  
          } else {  
            return n * factorial(n-1);  
          }  
        }  
      }  
    }  
  }  
}
```

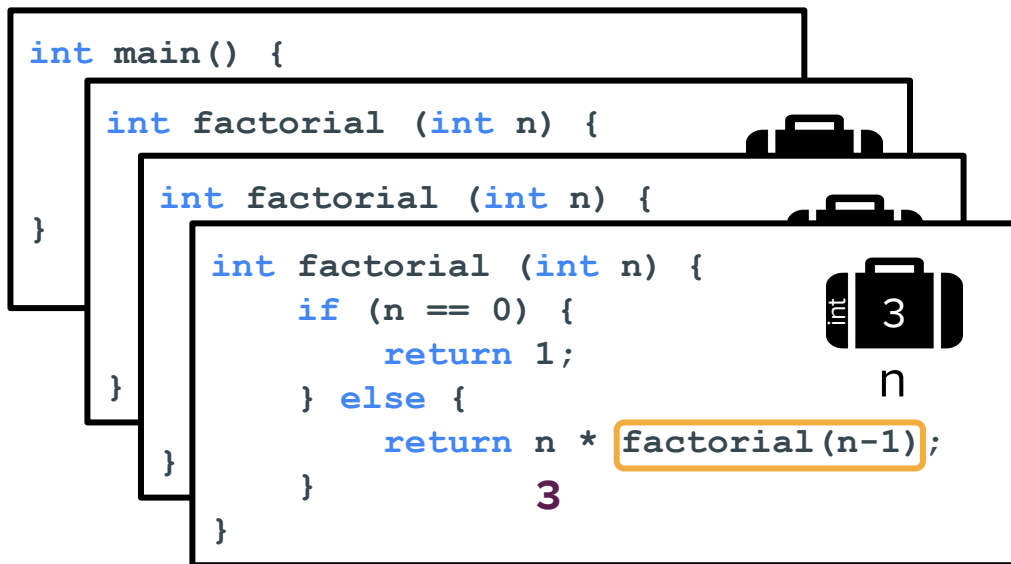
int 2

2 x 1

Recursion in action



Recursion in action

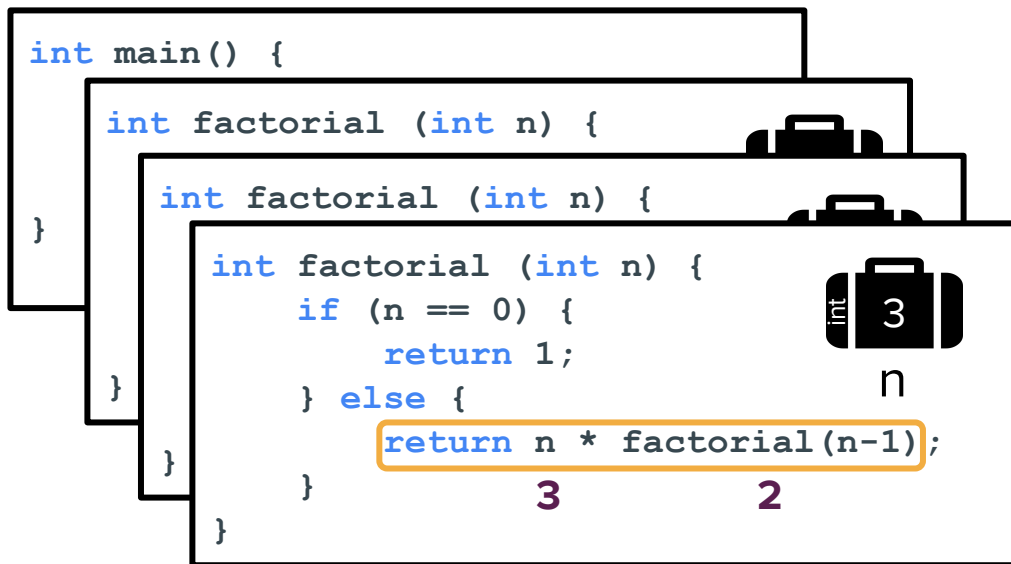


Recursion in action

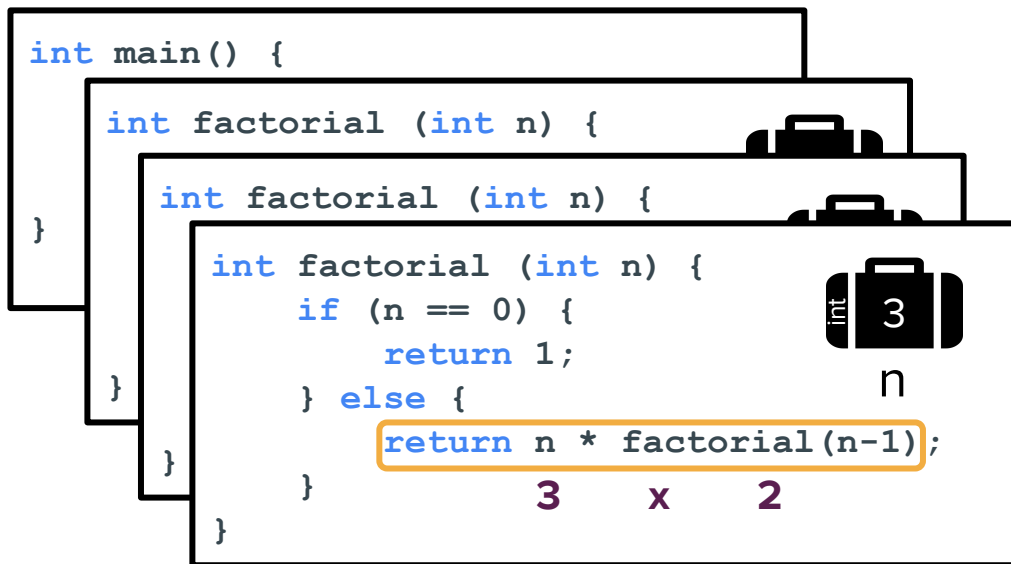
```
int main() {  
  int factorial (int n) {  
    int factorial (int n) {  
      int factorial (int n) {  
        if (n == 0) {  
          return 1;  
        } else {  
          return n * factorial(n-1);  
        }  
      }  
    }  
  }  
}
```

The diagram illustrates the recursive process for calculating the factorial of 3. It shows four overlapping boxes representing function frames. The innermost box shows the call to `factorial(2)` from within `factorial(3)`. A suitcase icon labeled `int 3` is positioned above the `n` parameter in the innermost box. The `return` statement `return n * factorial(n-1);` is highlighted with an orange box, with `3` and `2` written below `n` and `n-1` respectively.

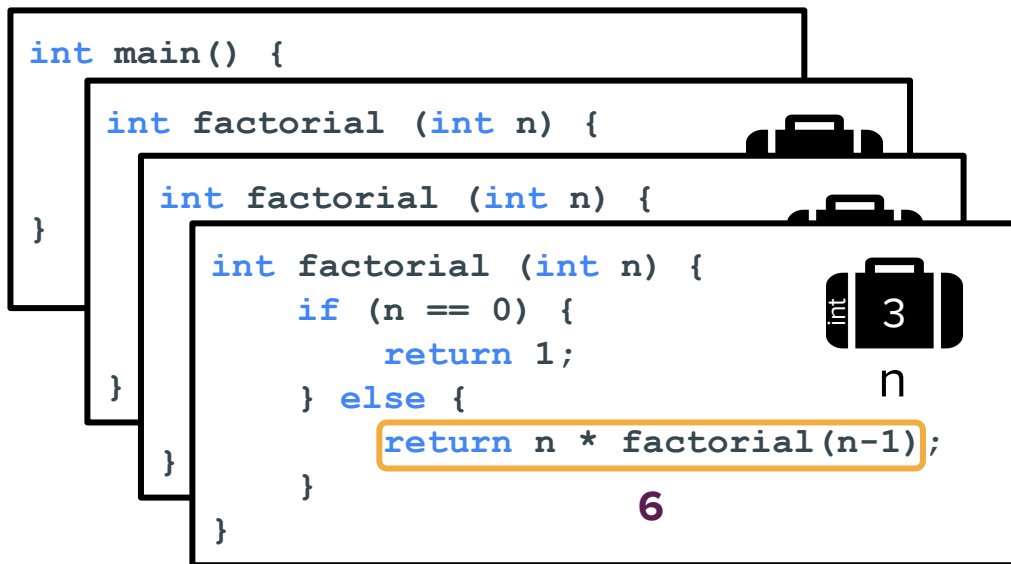
Recursion in action



Recursion in action



Recursion in action



Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
        int factorial (int n) {
```

```
            if (n == 0) {
```

```
                return 1;
```

```
            } else {
```

```
                return n * factorial(n-1);
```

```
            }
```

```
        }
```



n

4

factorial(n-1)

Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
        int factorial (int n) {
```

```
            if (n == 0) {
```

```
                return 1;
```

```
            } else {
```

```
                return n * factorial(n-1);
```

```
            }
```

```
        }
```



n

4

6

Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
        int factorial (int n) {
```

```
            if (n == 0) {
```

```
                return 1;
```

```
            } else {
```

```
                return n * factorial(n-1);
```

4

6



n

Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
        int factorial (int n) {
```

```
            if (n == 0) {
```

```
                return 1;
```

```
            } else {
```

```
                return n * factorial(n-1);
```

```
            }
```

```
        }
```



n

4

x

6

Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
        int factorial (int n) {
```

```
            if (n == 0) {
```

```
                return 1;
```

```
            } else {
```

```
                return n * factorial(n-1);
```

24



n

Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
        if (n == 0) {
```

```
            return 1;
```

```
        } else {
```

```
            return n * factorial(n-1);
```

```
        }
```

```
    }
```



`n`

5

Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
        if (n == 0) {
```

```
            return 1;
```

```
        } else {
```

```
            return n * factorial(n-1);
```

```
        }
```

```
    }
```



n

5

24

Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
        if (n == 0) {
```

```
            return 1;
```

```
        } else {
```

```
            return n * factorial(n-1);
```

```
        }
```

```
    }
```



n

5

24

Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
        if (n == 0) {
```

```
            return 1;
```

```
        } else {
```

```
            return n * factorial(n-1);
```

```
        }
```

```
    }
```



n

5 x **24**

Recursion in action

```
int main() {
```

```
    int factorial (int n) {
```

```
        if (n == 0) {
```

```
            return 1;
```

```
        } else {
```

```
            return n * factorial(n-1);
```

```
        }
```

```
    }
```



n

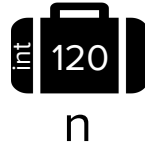
120

Recursion in action

```
int main() {  
    int n = factorial(5);  
    cout << "5! = " << n << endl;  
    return 0;  
}
```

Recursion in action

```
int main() {  
    int n = factorial(5);  
    cout << "5! = " << n << endl;  
    return 0;  
}
```



int 120
n

Recursive vs. Iterative

[Qt Creator]

Reverse string example

How can we reverse a string?

Suppose we want to reverse strings like in the following examples:

“dog” → “god”

“stressed” → “desserts”

“recursion” → “noisrucer”

“level” → “level”

“a” → “a”

Approaching recursive problems

- Look for self-similarity.
- Try out an example.
 - Work through a simple example and then increase the complexity.
 - Think about what information needs to be “stored” at each step in the recursive case (like the current value of **n** in each **factorial** stack frame).
- Ask yourself:
 - What is the base case? (What is the simplest case?)
 - What is the recursive case? (What pattern of self-similarity do you see?)

Discuss:

What are the base and
recursive cases?

(breakout rooms)

How can we reverse a string?

- Look for self-similarity: **stressed** → **desserts**

How can we reverse a string?

- Look for self-similarity: **stressed** → **desserts**
 - What's the first step you would take to reverse “stressed”?

How can we reverse a string?

- Look for self-similarity: **stressed** → **desserts**
 - Take the s and put it at the end of the string.

How can we reverse a string?

- Look for self-similarity: **stressed** → **desserts**
 - Take the s and put it at the end of the string.
 - Then reverse “tressed”

How can we reverse a string?

- Look for self-similarity: **stressed** → **desserts**
 - Take the s and put it at the end of the string.
 - Then reverse “tressed”:
 - Take the t and put it at the end of the string.
 - Then reverse “ressed”

How can we reverse a string?

- Look for self-similarity: **stressed** → **desserts**
 - Take the s and put it at the end of the string.
 - Then reverse “tressed”:
 - Take the t and put it at the end of the string.
 - Then reverse “ressed”:
 - Take the r and put it at the end of the string.
 - Then reverse “essed”

How can we reverse a string?

- Look for self-similarity: **stressed** → **desserts**
 - Take the s and put it at the end of the string.
 - Then reverse “tressed”:
 - Take the t and put it at the end of the string.
 - Then reverse “ressed”:
 - Take the r and put it at the end of the string.
 - Then reverse “essed”:
 - ...
 - Take the d and put it at the end of the string.
 - **Base case:** reverse “” → get “”

How can we reverse a string?

*How can we
express the
recursive case?*

- Look for self-similarity: **stressed** → **desserts**
 - Take the s and put it at the end of the string.
 - Then reverse “tressed”:
 - Take the t and put it at the end of the string.
 - Then reverse “ressed”:
 - Take the r and put it at the end of the string.
 - Then reverse “essed”:
 - ...
 - Take the d and put it at the end of the string.
 - **Base case:** reverse “” → get “”

How can we reverse a string?

*How can we
express the
recursive case?*

- Look for self-similarity: **stressed** → **desserts**
 - **Take the s and put it at the end of the string.**
 - **Then reverse “tressed”:**
 - Take the t and put it at the end of the string.
 - Then reverse “ressed”:
 - Take the r and put it at the end of the string.
 - Then reverse “essed”:
 - ...
 - Take the d and put it at the end of the string.
 - **Base case:** reverse “” → get “”

How can we reverse a string?

*How can we
express the
recursive case?*

- Look for self-similarity: **stressed** → **desserts**
 - **reverse("stressed") = reverse("tressed") + 's'**
 - Take the t and put it at the end of the string.
 - Then reverse "ressed":
 - Take the r and put it at the end of the string.
 - Then reverse "essed":
 - ...
 - Take the d and put it at the end of the string.
 - **Base case:** reverse "" → get ""

How can we reverse a string?

*How can we
express the
recursive case?*

- Look for self-similarity: **stressed** → **desserts**
 - reverse(“stressed”) = reverse(“tressed”) + ‘s’
 - **Take the t and put it at the end of the string.**
 - **Then reverse “ressed”:**
 - Take the r and put it at the end of the string.
 - Then reverse “essed”:
 - ...
 - Take the d and put it at the end of the string.
 - **Base case:** reverse “” → get “”

How can we reverse a string?

*How can we
express the
recursive case?*

- Look for self-similarity: **stressed** → **desserts**
 - reverse("stressed") = reverse("tressed") + 's'
 - **reverse("tressed") = reverse("ressed") + 't'**
 - Take the r and put it at the end of the string.
 - Then reverse "essed":
 - ...
 - Take the d and put it at the end of the string.
 - **Base case:** reverse "" → get ""

How can we reverse a string?

*How can we
express the
recursive case?*

- Look for self-similarity: **stressed** → **desserts**
 - reverse(“stressed”) = reverse(“tressed”) + ‘s’
 - reverse(“tressed”) = reverse(“ressed”) + ‘t’
 - **Take the r and put it at the end of the string.**
 - **Then reverse “essed”:**
 - ...
 - Take the d and put it at the end of the string.
 - **Base case:** reverse “” → get “”

How can we reverse a string?

*How can we
express the
recursive case?*

- Look for self-similarity: **stressed** → **desserts**
 - $\text{reverse}(\text{"stressed"}) = \text{reverse}(\text{"tressed"}) + \text{'s'}$
 - $\text{reverse}(\text{"tressed"}) = \text{reverse}(\text{"ressed"}) + \text{'t'}$
 - $\text{reverse}(\text{"ressed"}) = \text{reverse}(\text{"essed"}) + \text{'r'}$
 - ...
 - Take the d and put it at the end of the string.
 - **Base case:** reverse "" → get ""

How can we reverse a string?

- Look for self-similarity: **stressed** → **desserts**
 - $\text{reverse}(\text{"stressed"}) = \text{reverse}(\text{"tressed"}) + \text{'s'}$
 - $\text{reverse}(\text{"tressed"}) = \text{reverse}(\text{"ressed"}) + \text{'t'}$
 - $\text{reverse}(\text{"ressed"}) = \text{reverse}(\text{"essed"}) + \text{'r'}$
 - ...
 - **Base case:** $\text{reverse}(\text{""}) = \text{""}$

How can we reverse a string?

- **Recursive case:** $\text{reverse}(\text{str}) = \text{reverse}(\text{str without first letter}) + \text{first letter of str}$
- **Base case:** $\text{reverse}("") = ""$

How can we reverse a string?

- **Recursive case:** $\text{reverse}(\text{str}) = \text{reverse}(\text{str without first letter}) + \text{first letter of str}$
- **Base case:** $\text{reverse}("") = ""$

Depending on how you thought of the problem, you may have also come up with:

- **Recursive case:** $\text{reverse}(\text{str}) = \text{last letter of str} + \text{reverse}(\text{str without last letter})$
- **Base case:** $\text{reverse}("") = ""$

Let's code it!

(live coding)

Summary

Summary

- Recursion is a problem-solving technique in which tasks are completed by reducing them into **repeated, smaller tasks of the same form**.
 - A recursive operation (function) is defined in terms of itself (i.e. it calls itself).

Summary

- Recursion is a problem-solving technique in which tasks are completed by reducing them into **repeated, smaller tasks of the same form**.
- Recursion has two main parts: the **base case** and the **recursive case**.
 - Base case: Simplest form of the problem that has a direct answer.
 - Recursive case: The step where you break the problem into a smaller, self-similar task.

Summary

- Recursion is a problem-solving technique in which tasks are completed by reducing them into **repeated, smaller tasks of the same form**.
- Recursion has two main parts: the **base case** and the **recursive case**.
- The solution will get built up **as you come back up the call stack**.
 - The base case will define the “base” of the solution you’re building up.
 - Each previous recursive call contributes a little bit to the final solution.
 - The initial call to your recursive function is what will return the completely constructed answer.

Summary

- Recursion is a problem-solving technique in which tasks are completed by reducing them into **repeated, smaller tasks of the same form**.
- Recursion has two main parts: the **base case** and the **recursive case**.
- The solution will get built up **as you come back up the call stack**.
- When solving problems recursively, look for **self-similarity** and think about **what information is getting stored in each stack frame**.

Summary

- Recursion is a problem-solving technique in which tasks are completed by reducing them into **repeated, smaller tasks of the same form**.
- Recursion has two main parts: the **base case** and the **recursive case**.
- The solution will get built up **as you come back up the call stack**.
- When solving problems recursively, look for **self-similarity** and think about **what information is getting stored in each stack frame**.

What's next?

Roadmap

C++ basics

User/client

vectors + grids

stacks + queues

sets + maps

Core
Tools

testing

algorithmic
analysis

recursive
problem-solving

Object-Oriented
Programming

Implementation

arrays

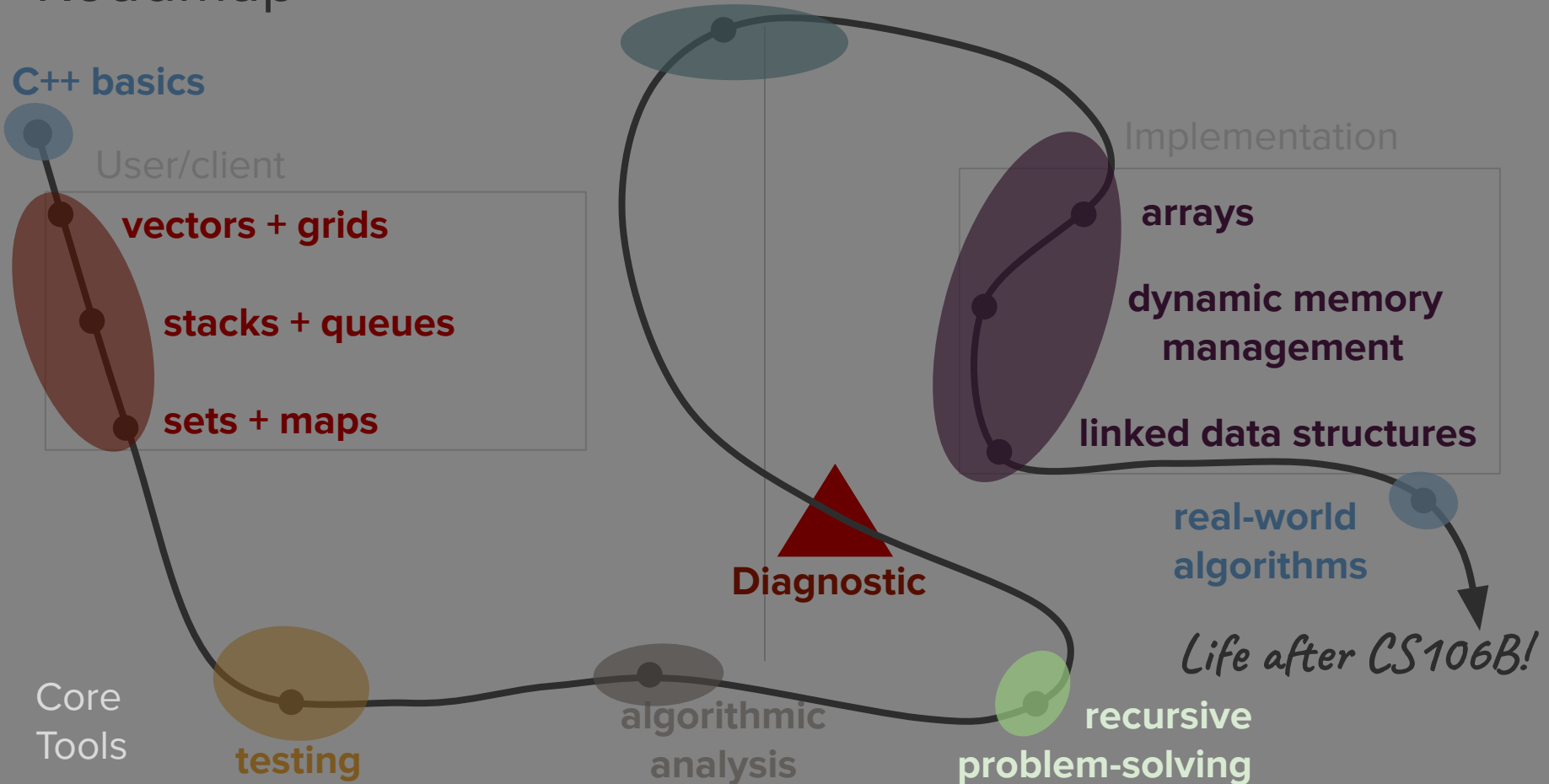
dynamic memory
management

linked data structures

real-world
algorithms

Life after CS106B!

Diagnostic



Fractals

