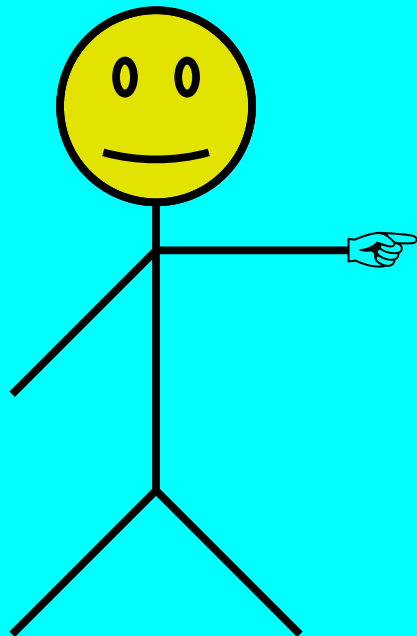


# Implementing Abstractions

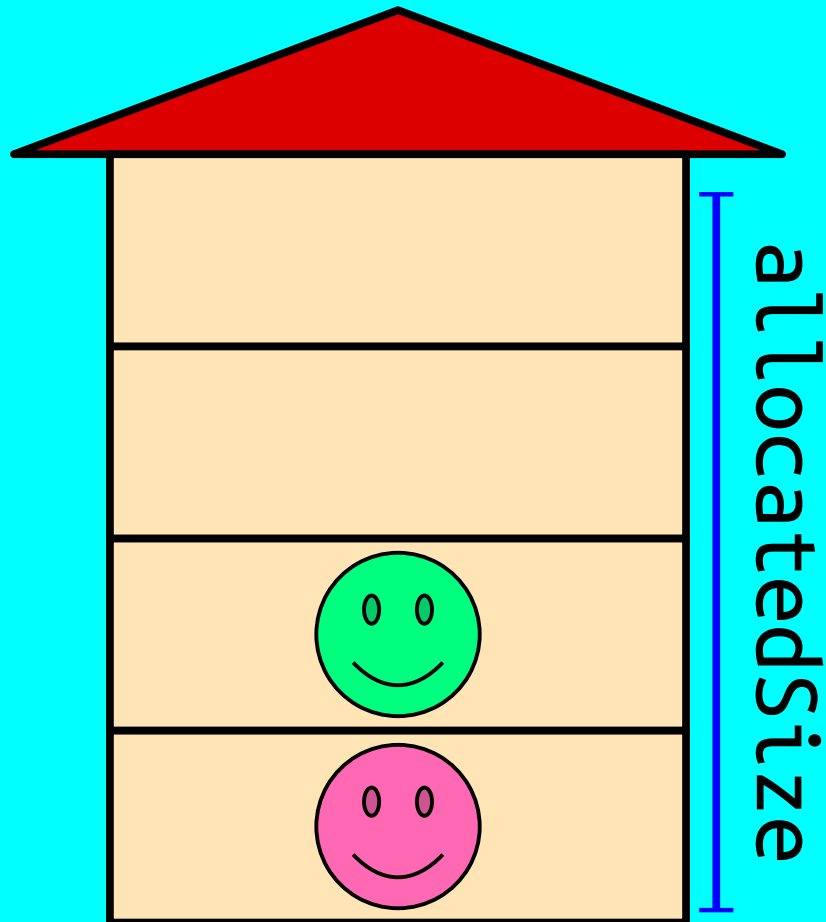
## Part Two

Previously, on CS106B...

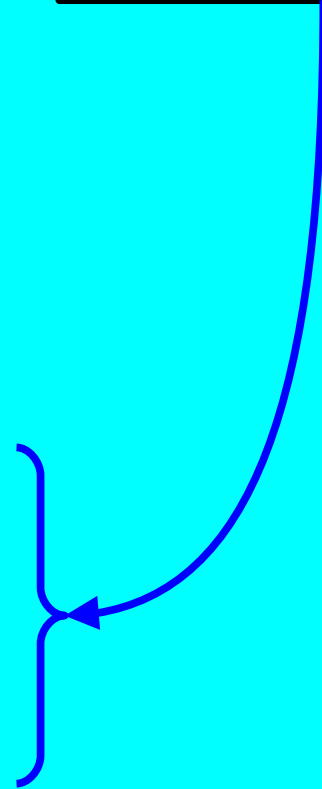
```
private:
    int* elems;
    int  allocatedSize;
    int  logicalSize;
```



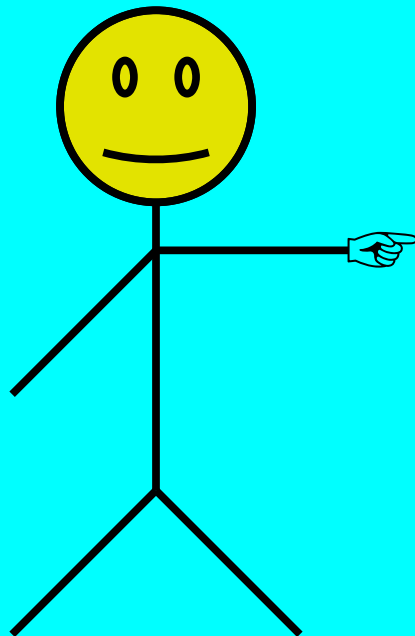
elems



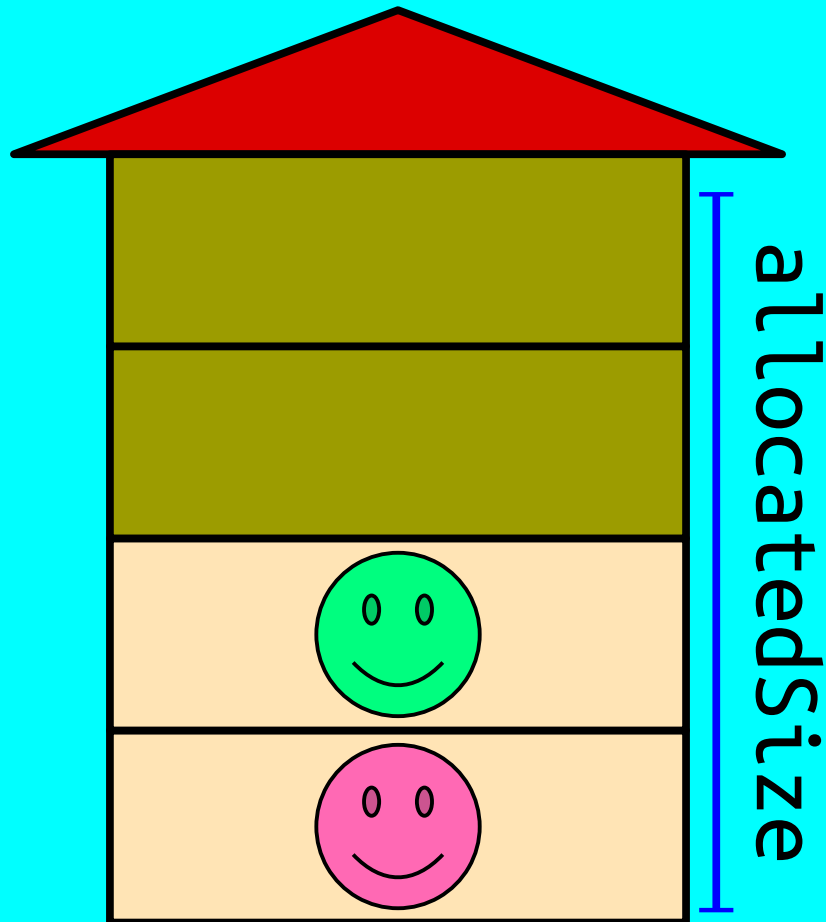
logicalSize



```
private:  
    int* elems;  
    int  allocatedSize;  
    int  logicalSize;
```



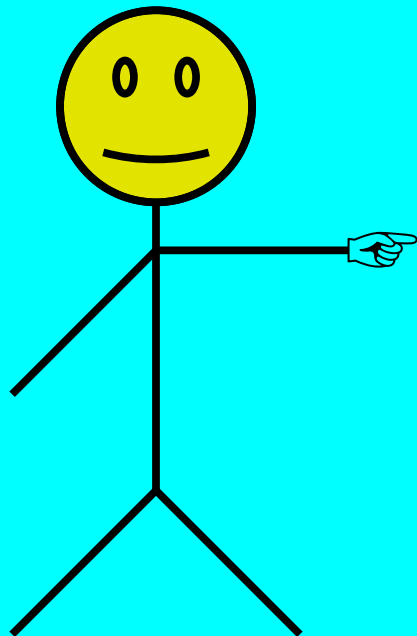
elems



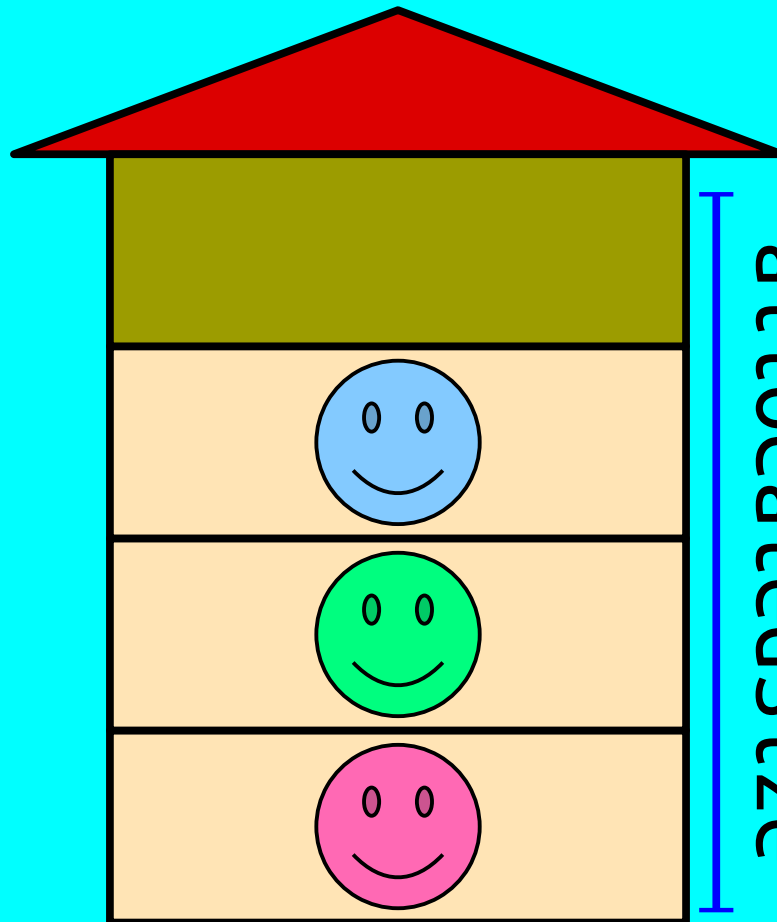
logicalSize



```
private:  
    int* elems;  
    int  allocatedSize;  
    int  logicalSize;
```

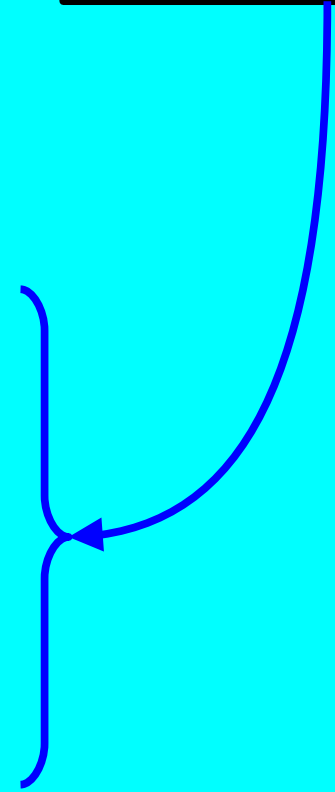


elems

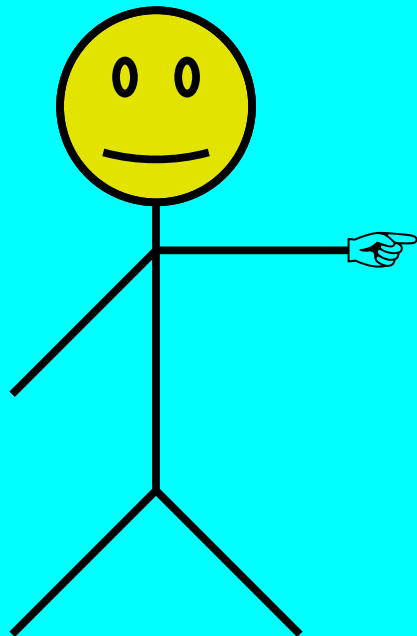


allocatedSize

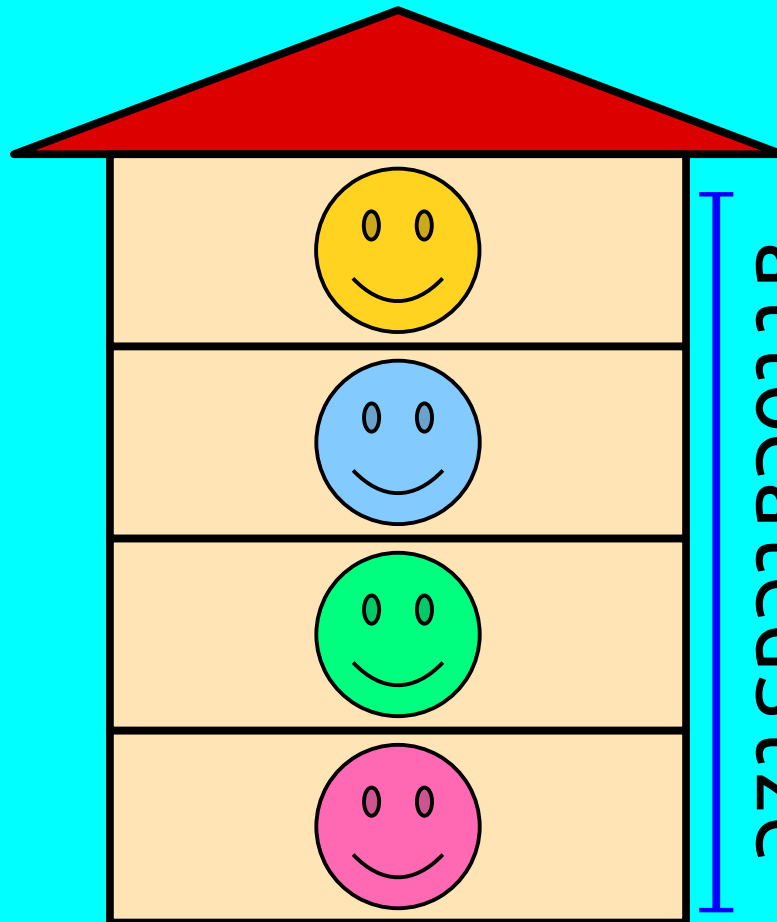
logicalSize



```
private:  
    int* elems;  
    int  allocatedSize;  
    int  logicalSize;
```

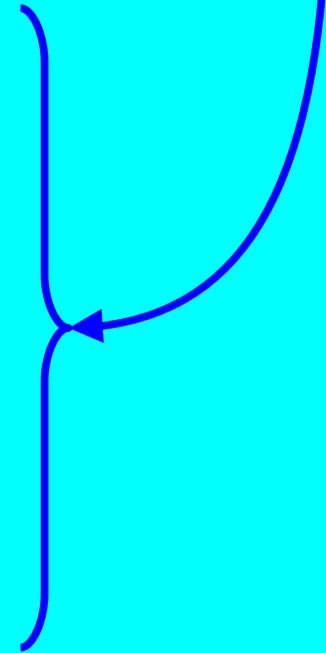


elems

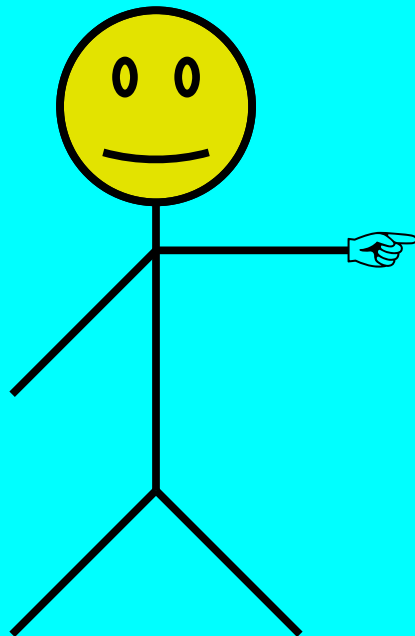


allocatedSize

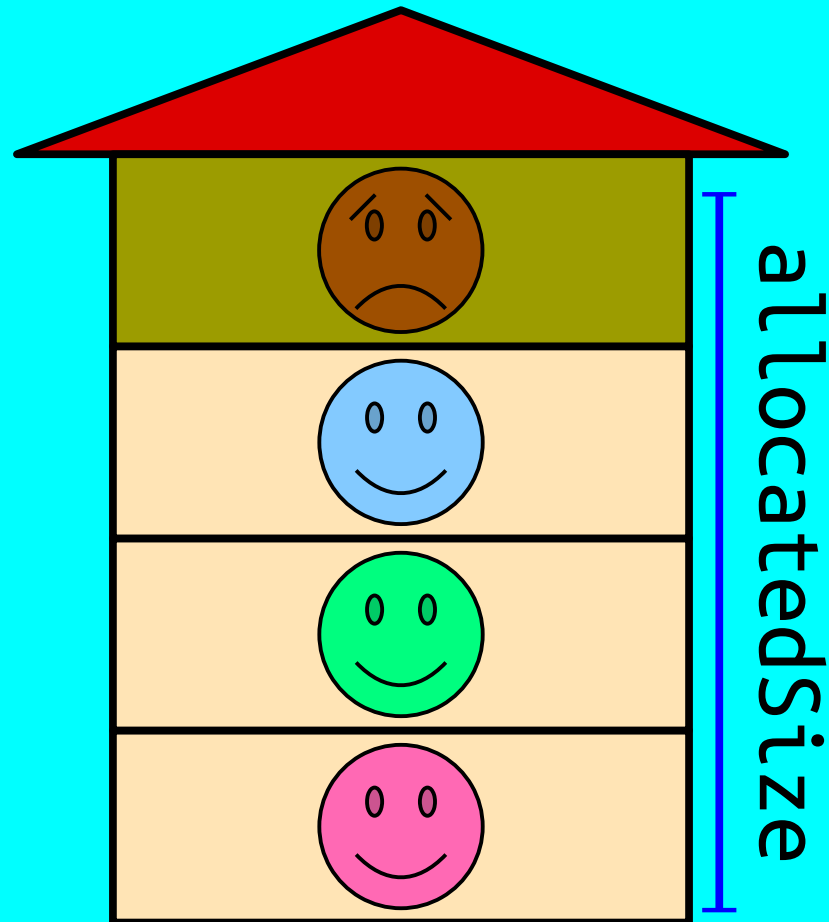
logicalSize



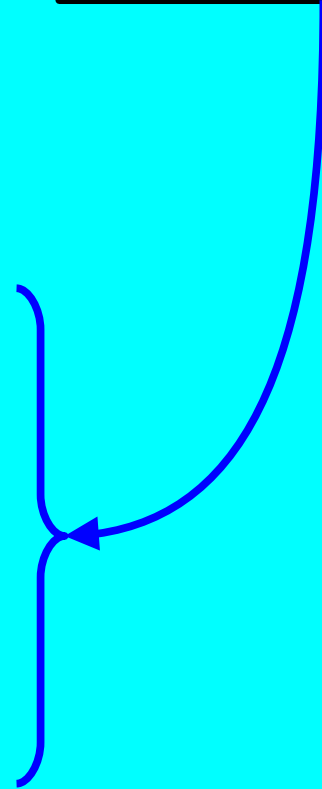
```
private:  
    int* elems;  
    int  allocatedSize;  
    int  logicalSize;
```



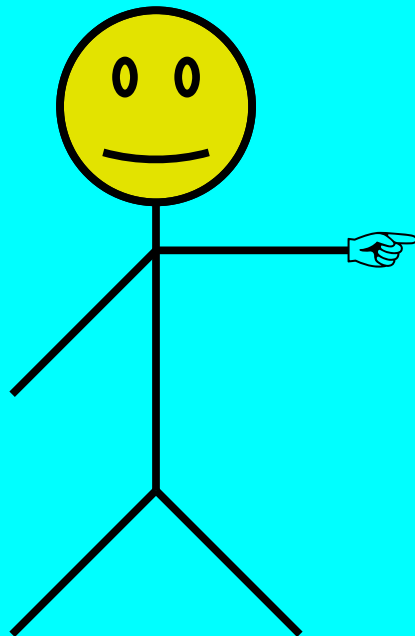
elems



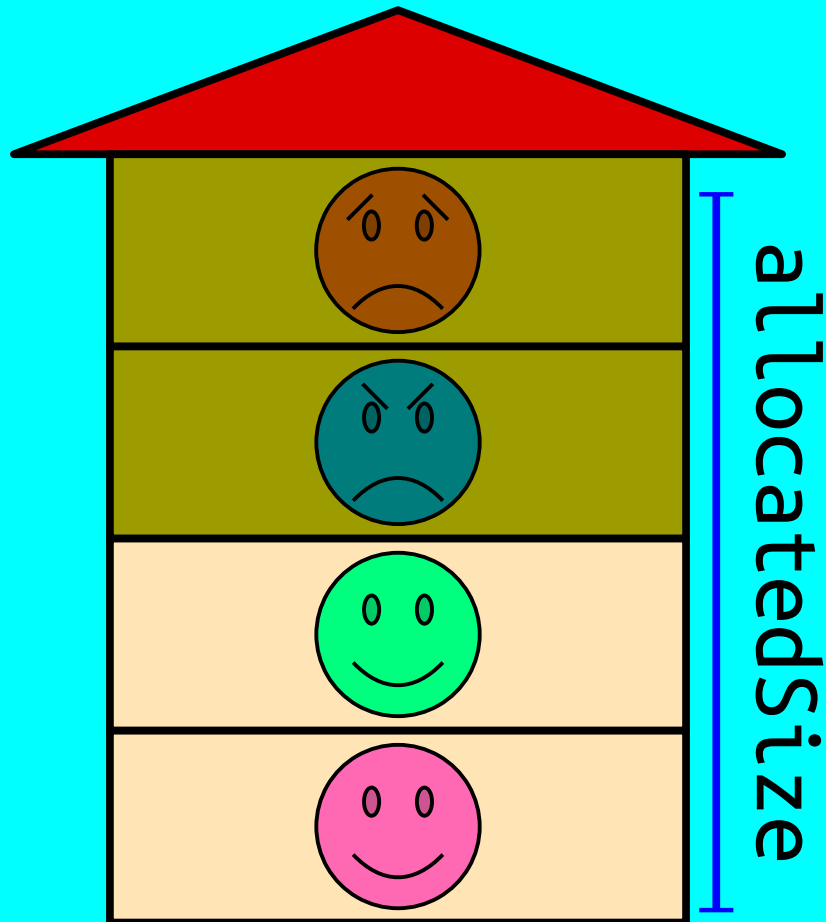
logicalSize



```
private:  
    int* elems;  
    int  allocatedSize;  
    int  logicalSize;
```

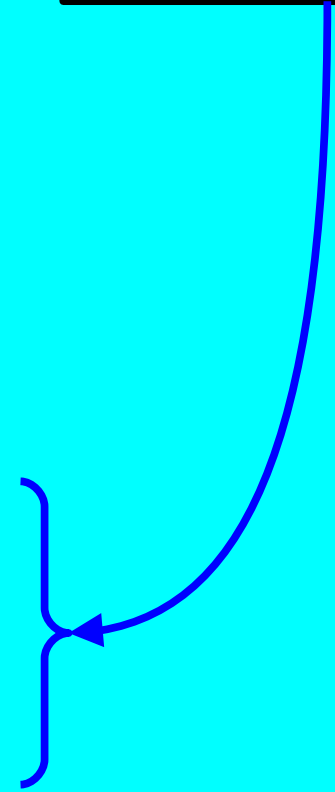


elems



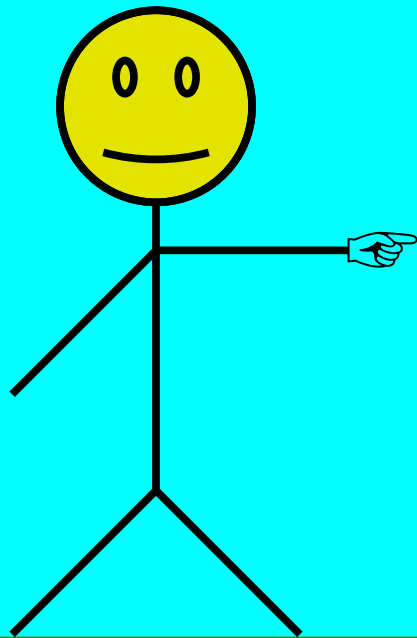
allocatedSize

logicalSize

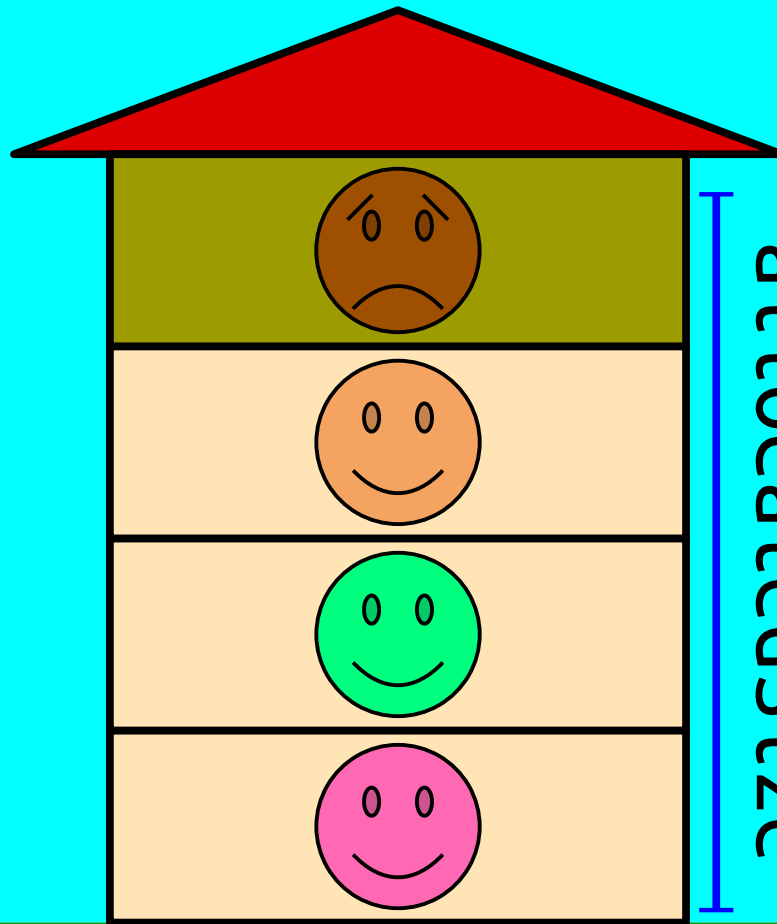




```
private:  
    int* elems;  
    int  allocatedSize;  
    int  logicalSize;
```

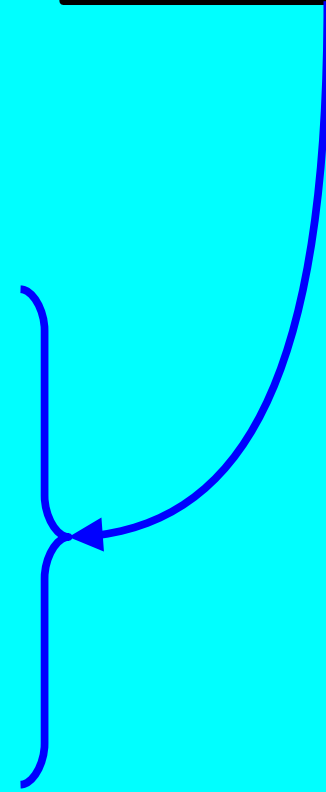


elems

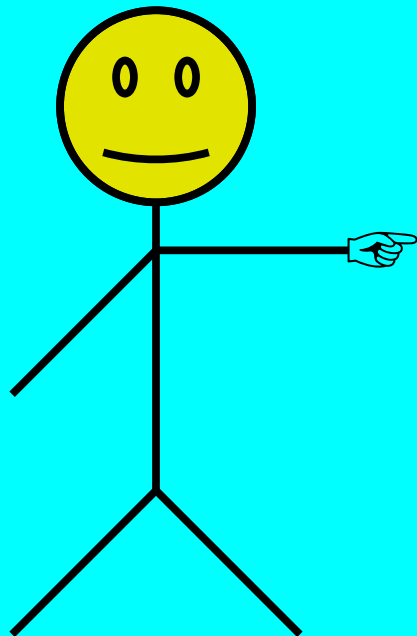


allocatedSize

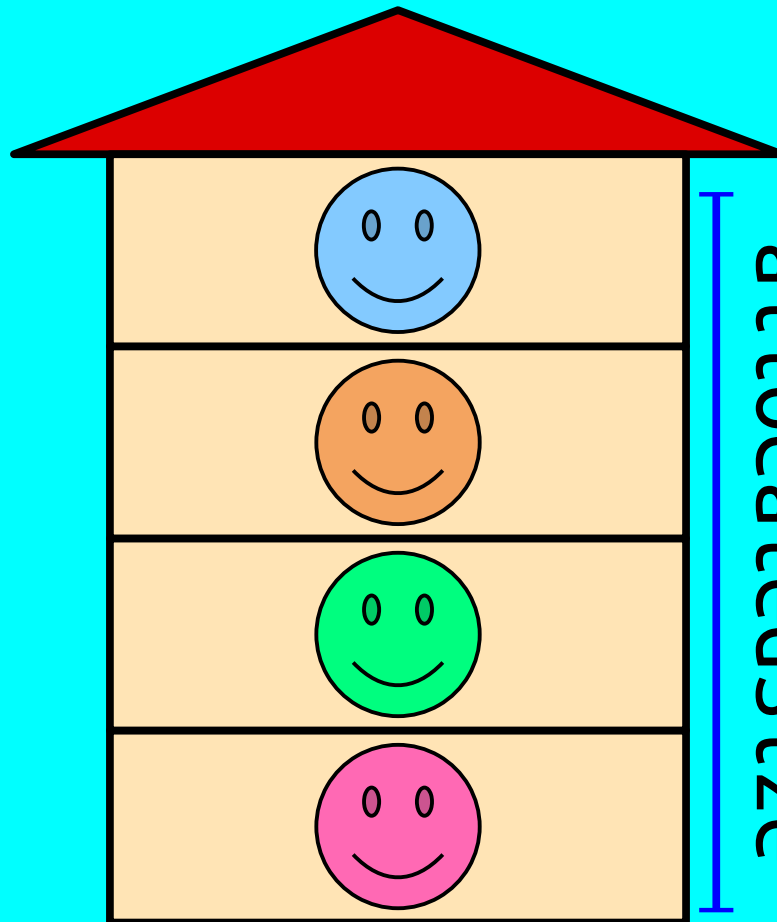
logicalSize



```
private:  
    int* elems;  
    int  allocatedSize;  
    int  logicalSize;
```

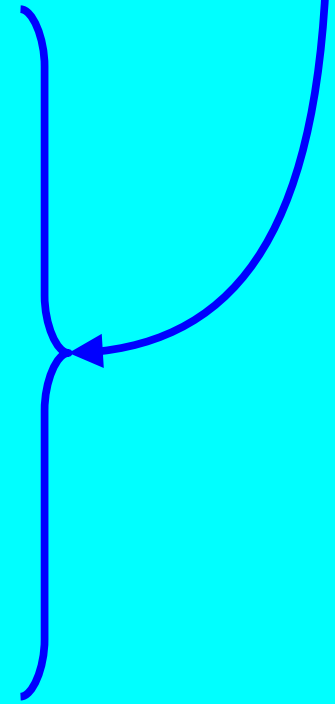


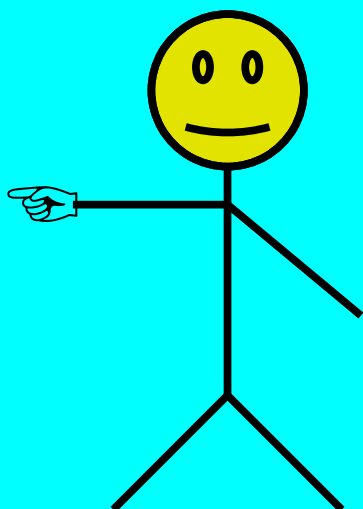
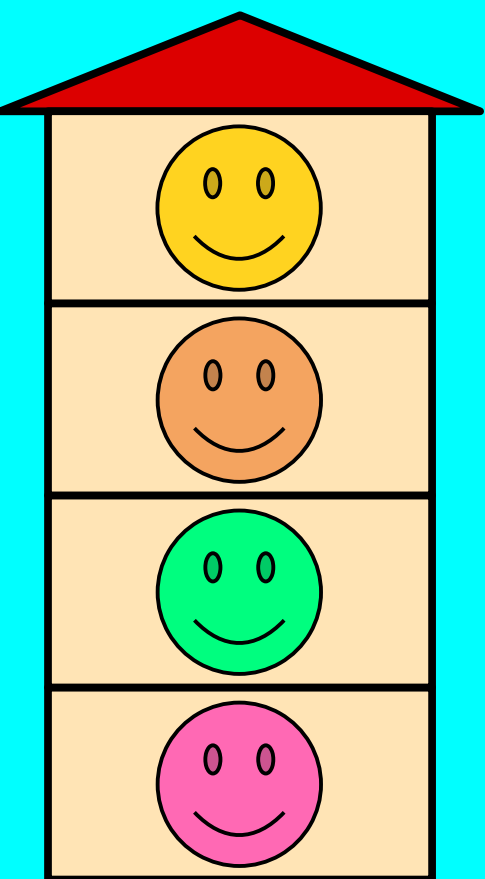
elems



allocatedSize

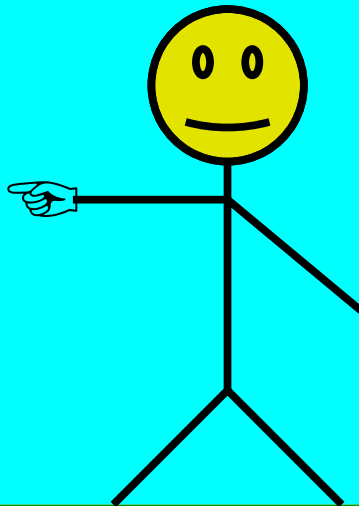
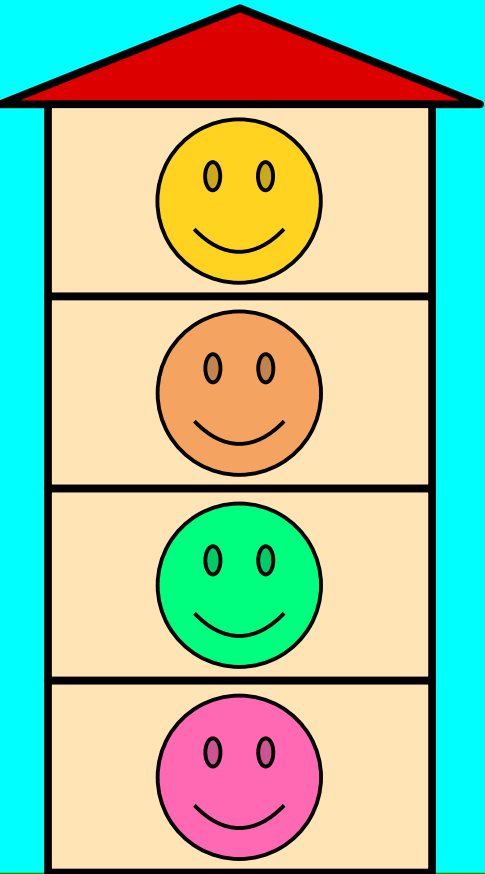
logicalSize



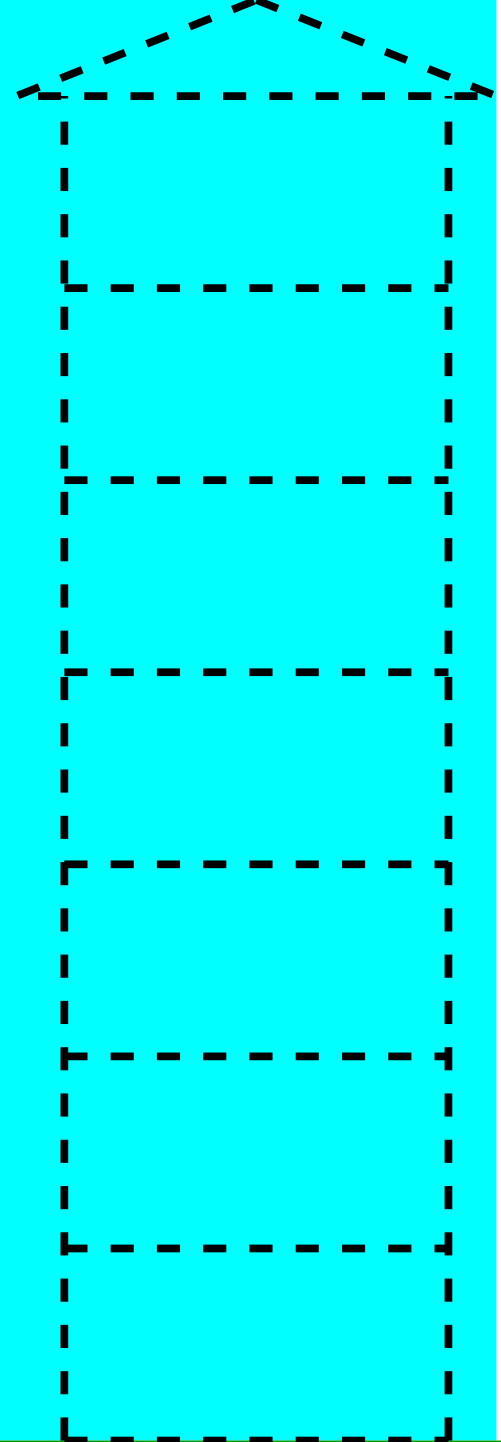


eLems

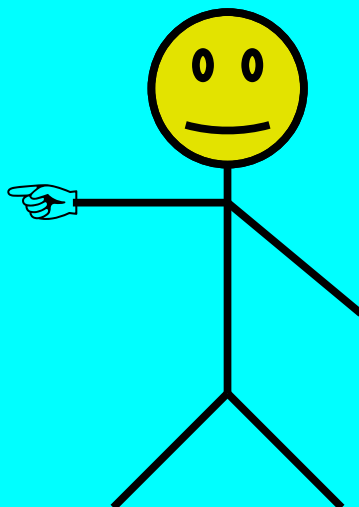
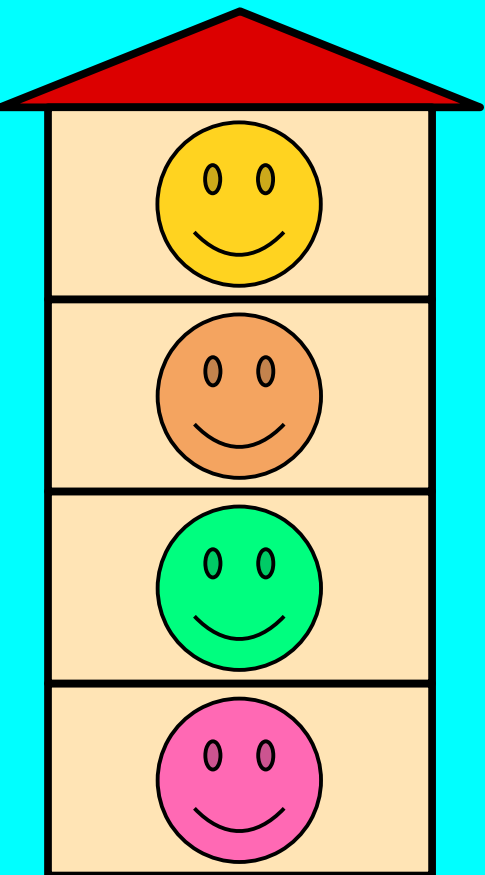
```
allocatedSize = /* bigger */;
```



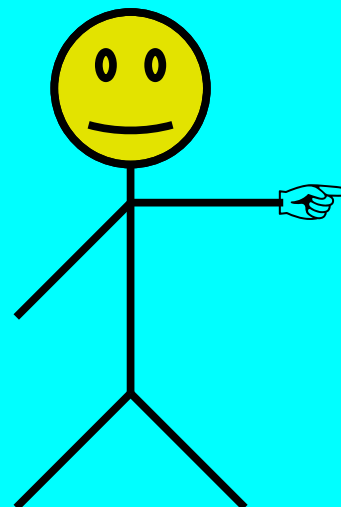
e1ems



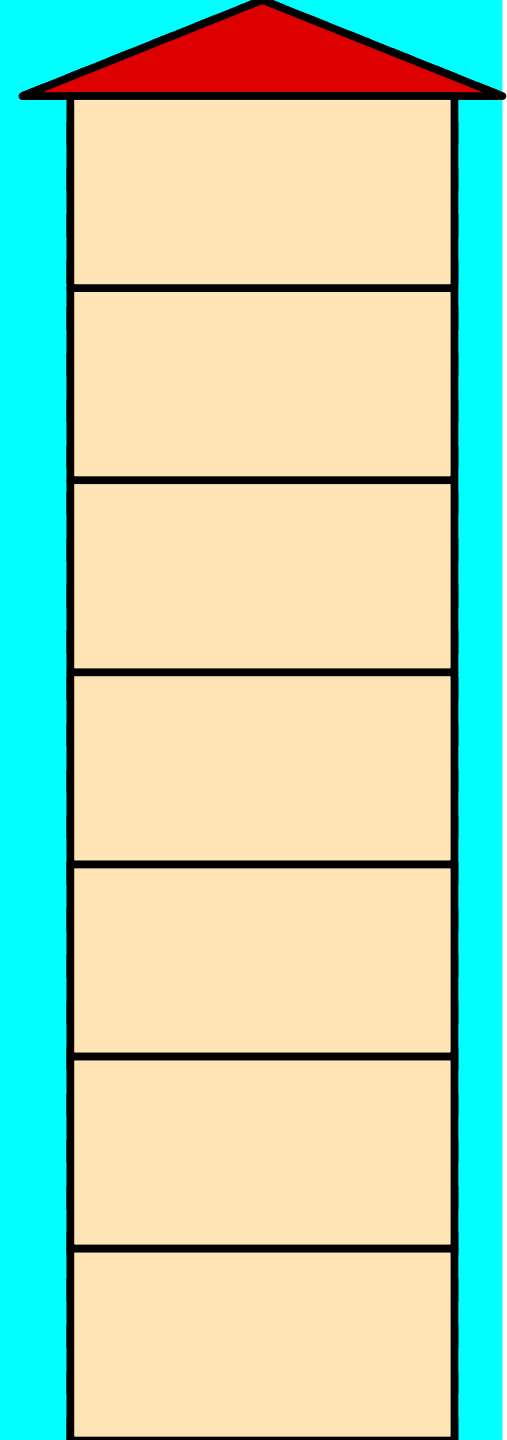
```
allocatedSize = /* bigger */;  
int* helper = new int[allocatedSize];
```



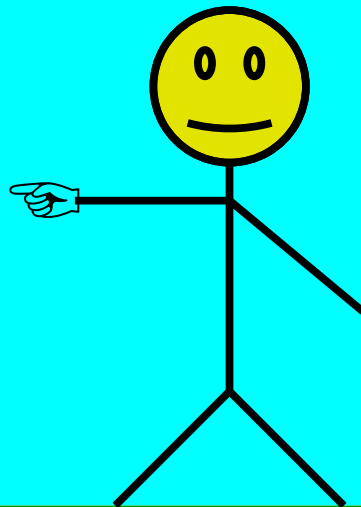
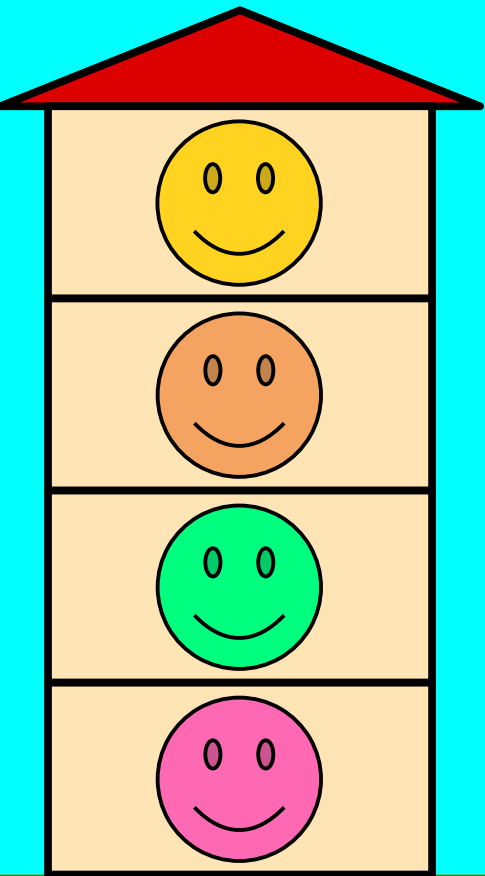
elems



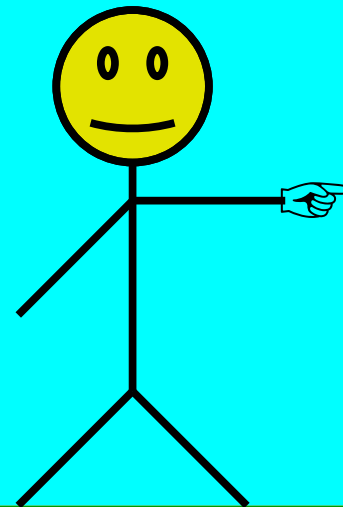
helper



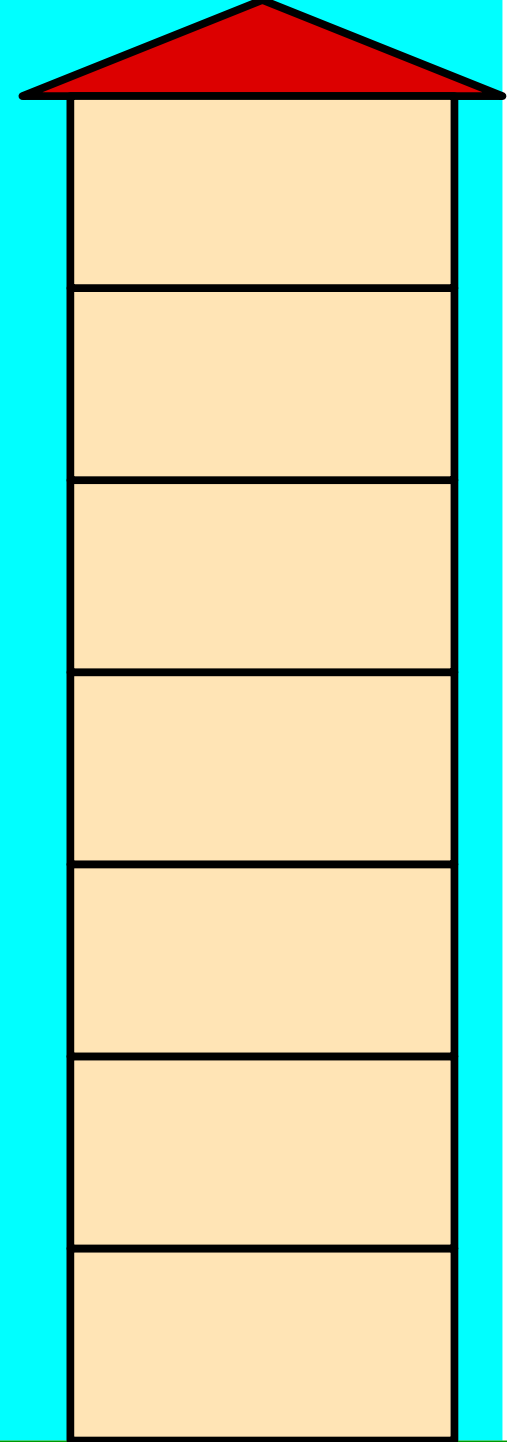
```
allocatedSize = /* bigger */;  
int* helper = new int[allocatedSize];  
/* ... move elements over ... */
```



elems

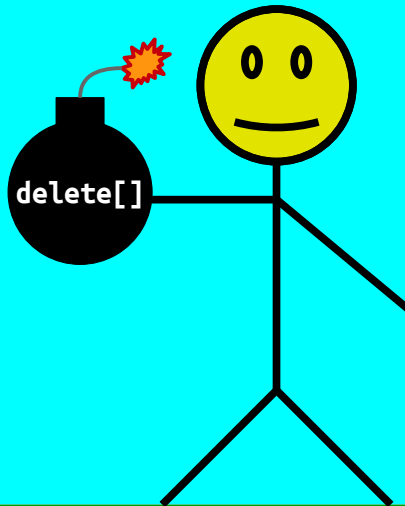


helper

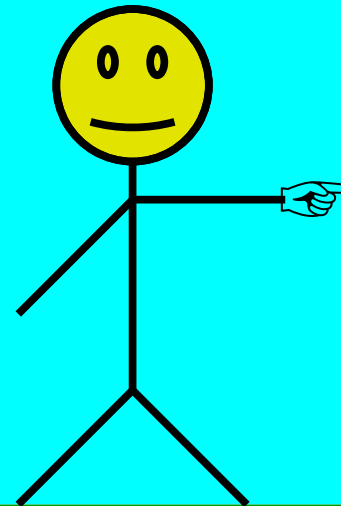


```
allocatedSize = /* bigger */;  
int* helper = new int[allocatedSize];  
/* ... move elements over ... */  
delete[] elems;
```

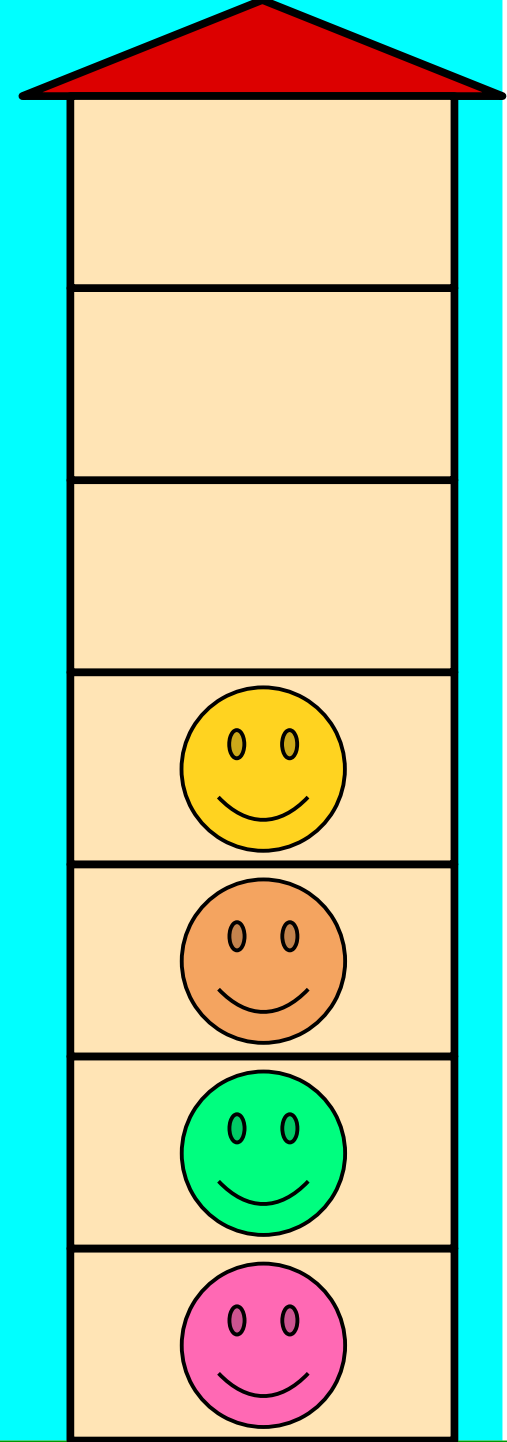
*Dynamic  
Deallocation!*



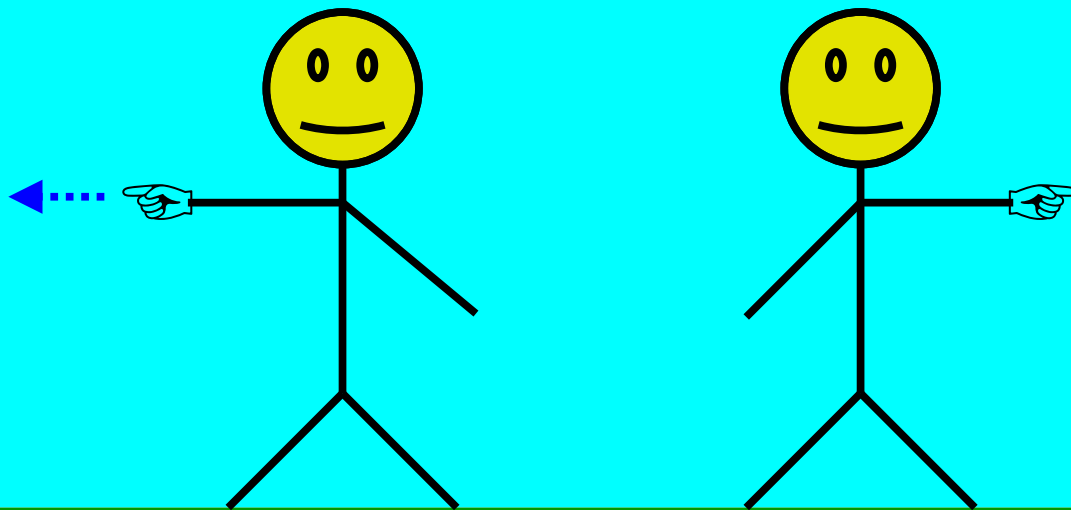
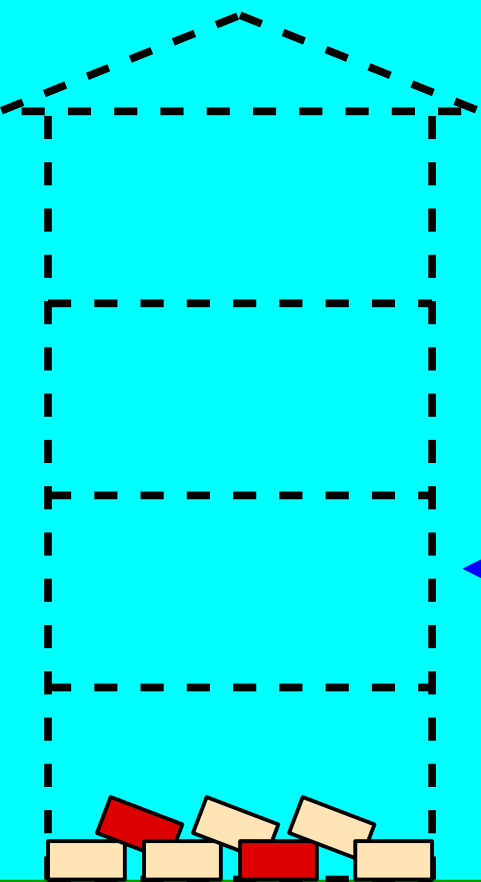
elems



helper

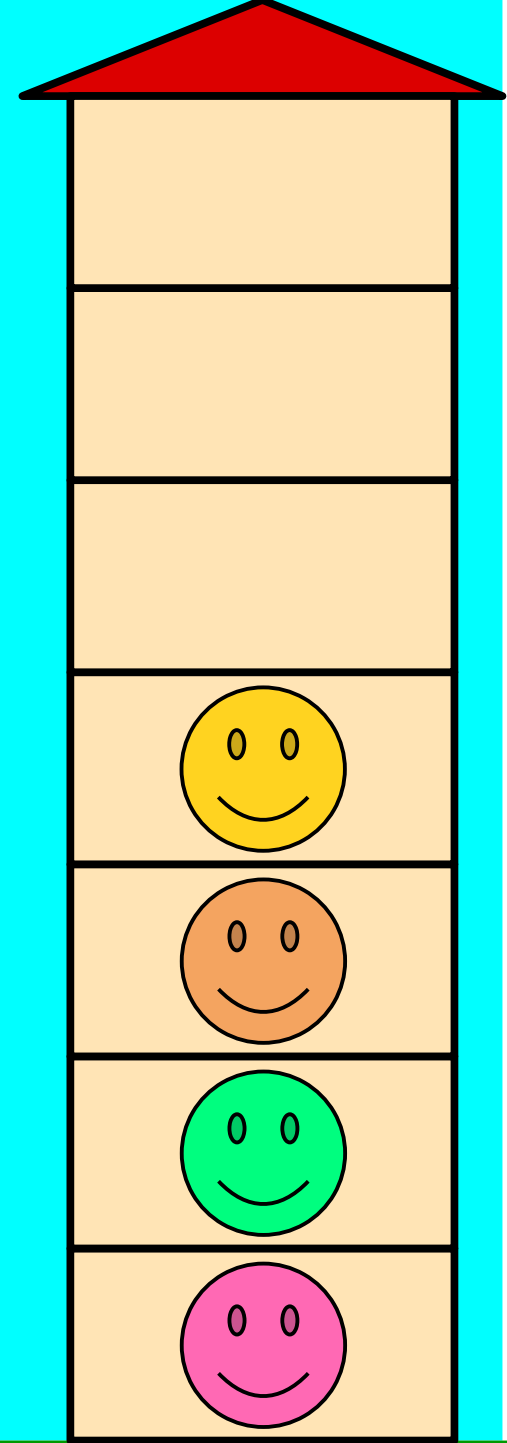


```
allocatedSize = /* bigger */;  
int* helper = new int[allocatedSize];  
/* ... move elements over ... */  
delete[] elems;
```



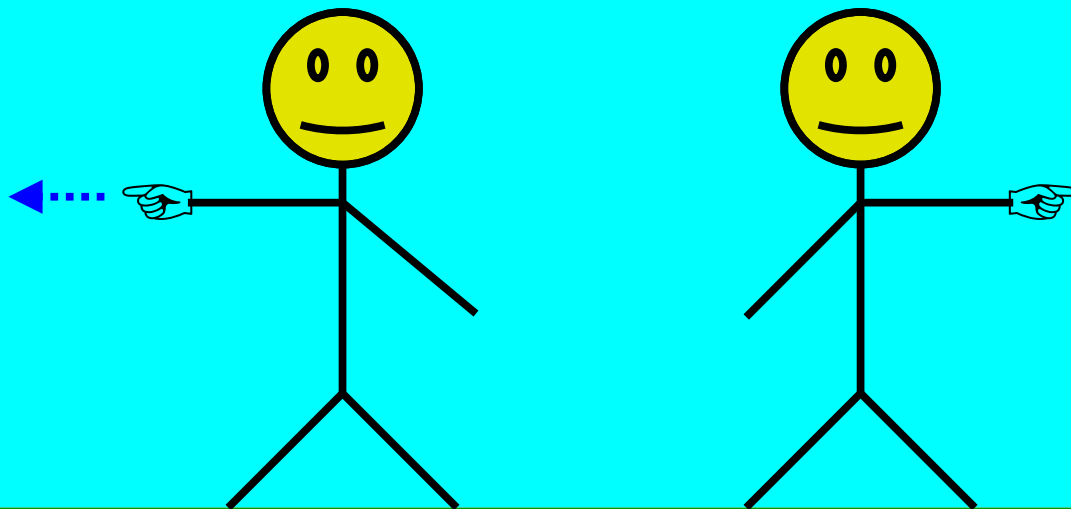
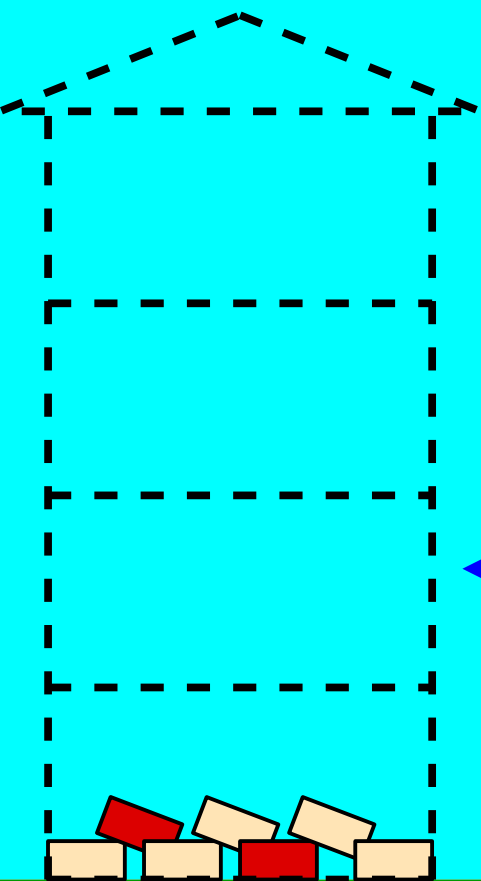
elems

helper



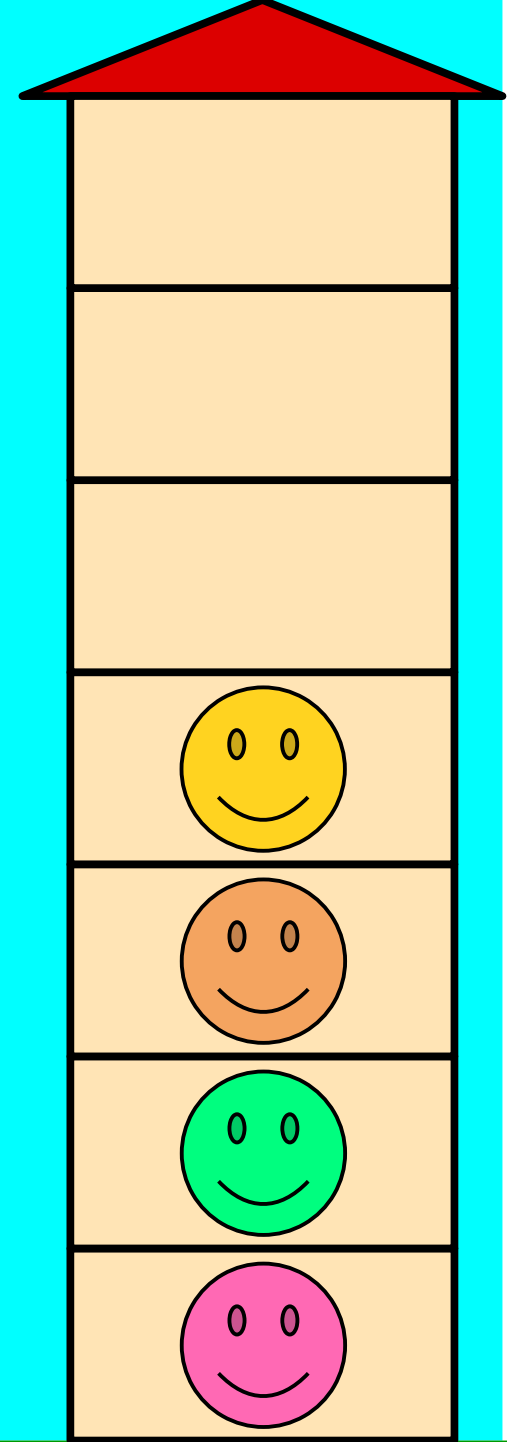


```
allocatedSize = /* bigger */;  
int* helper = new int[allocatedSize];  
/* ... move elements over ... */  
delete[] elems;  
elems = helper;
```

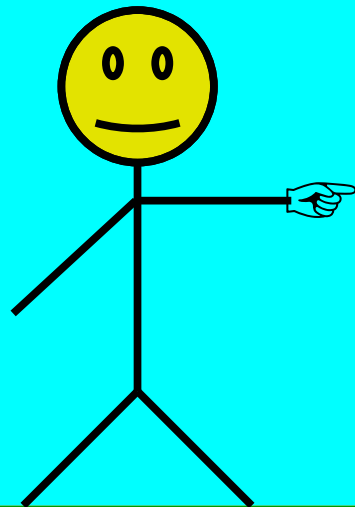
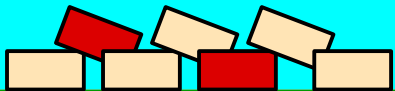


elems

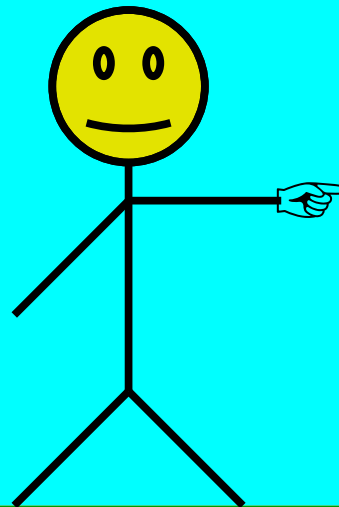
helper



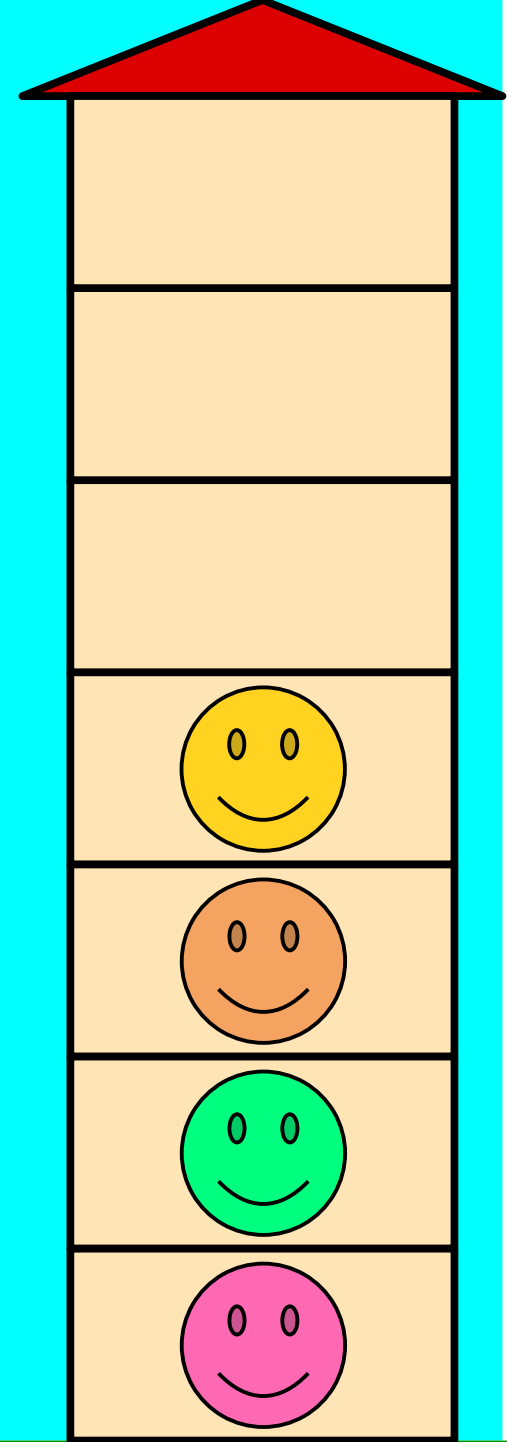
```
allocatedSize = /* bigger */;  
int* helper = new int[allocatedSize];  
/* ... move elements over ... */  
delete[] elems;  
elems = helper;
```



elems



helper



What is the big-O cost of a push?  
What is the big-O cost of  $n$  pushes?

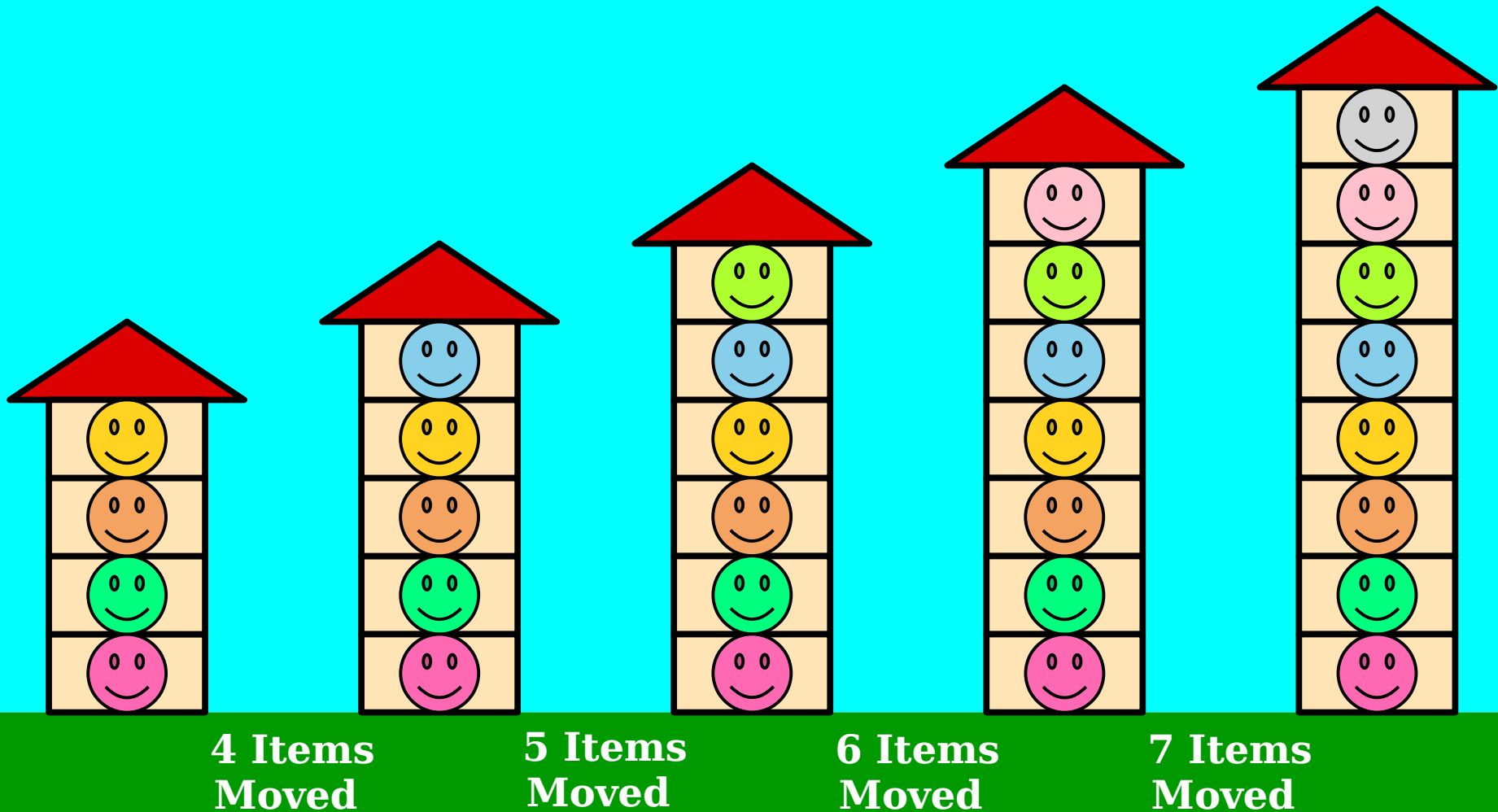
Take thirty seconds to think this over.  
Look over the runtime plot and the  
code for the push operation.

Formulate a hypothesis!

What is the big-O cost of a push?  
What is the big-O cost of  $n$  pushes?

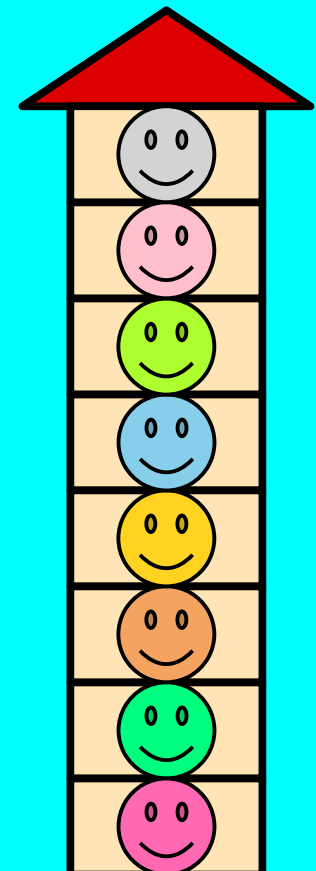
Turn to your neighbor  
and discuss!

What is the big-O cost of a push?  
What is the big-O cost of  $n$  pushes?



Every push beyond the first few requires moving all  $n$  elements from the old array to the new array.

Cost of a single push:  $O(n)$ .



4 Items  
Moved

5 Items  
Moved

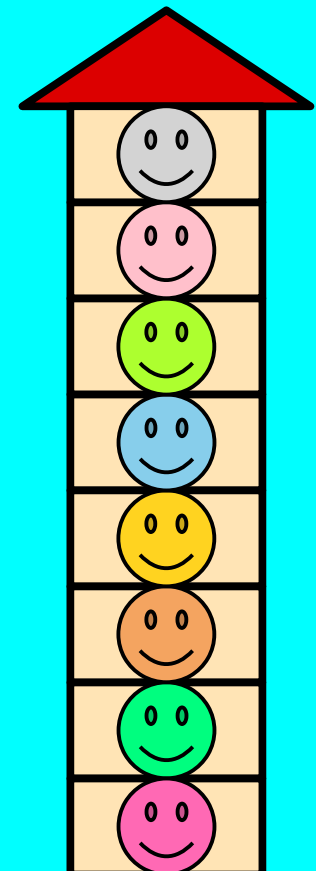
6 Items  
Moved

7 Items  
Moved

Every push beyond the first few requires moving all  $n$  elements from the old array to the new array.

Cost of doing  $n$  pushes:  
 $4 + 5 + 6 + \dots + n = \mathbf{O(n^2)}$ .

**Question:** How do we speed this up?

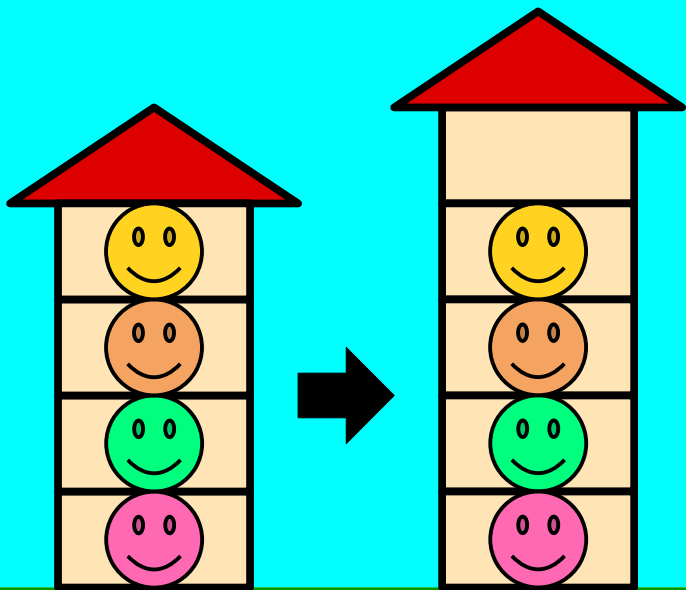


4 Items  
Moved

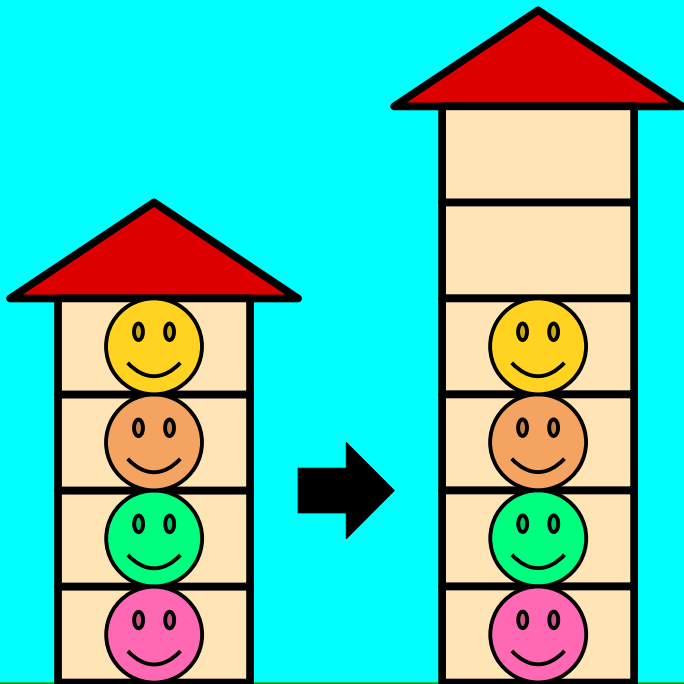
5 Items  
Moved

6 Items  
Moved

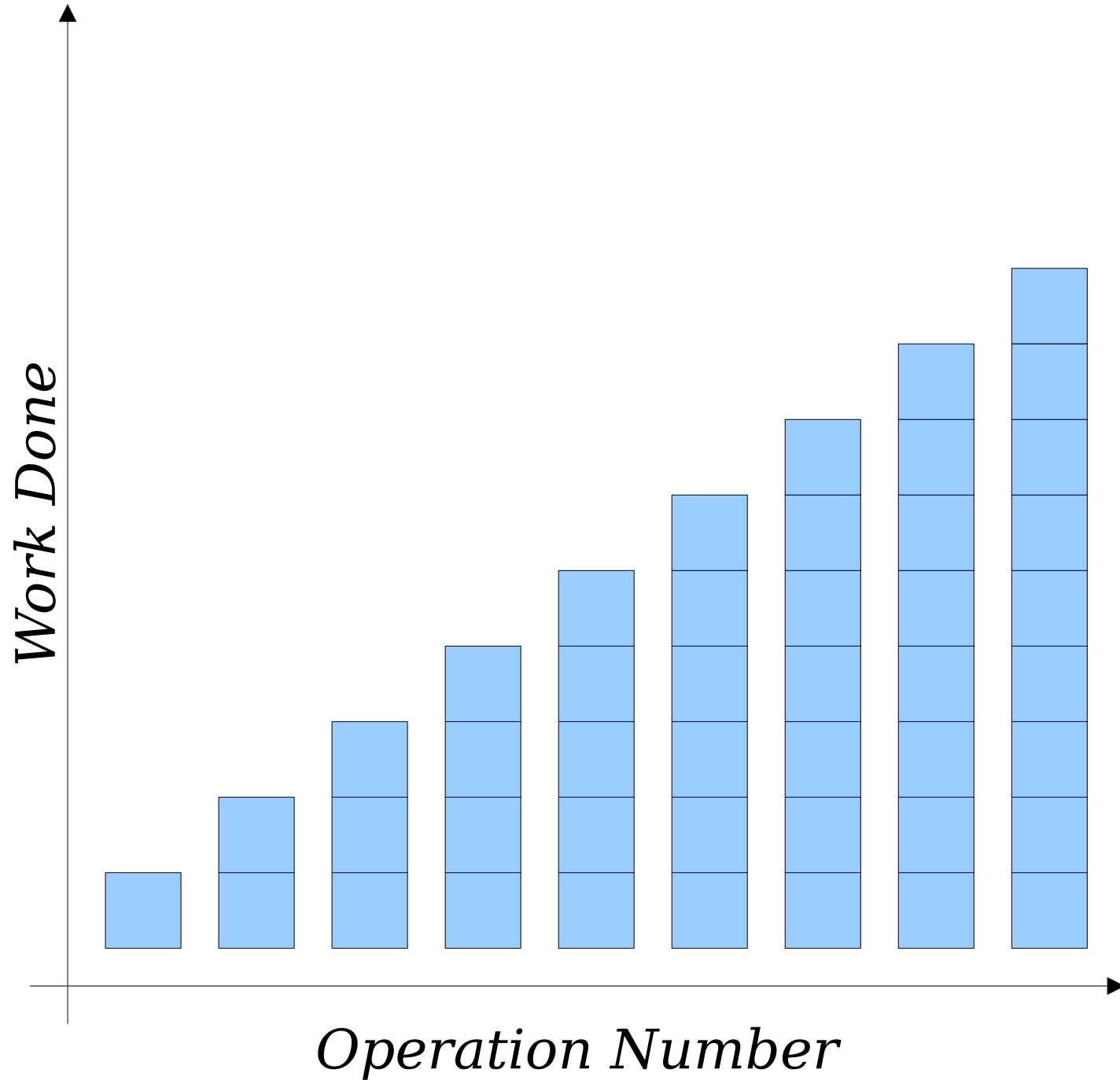
7 Items  
Moved



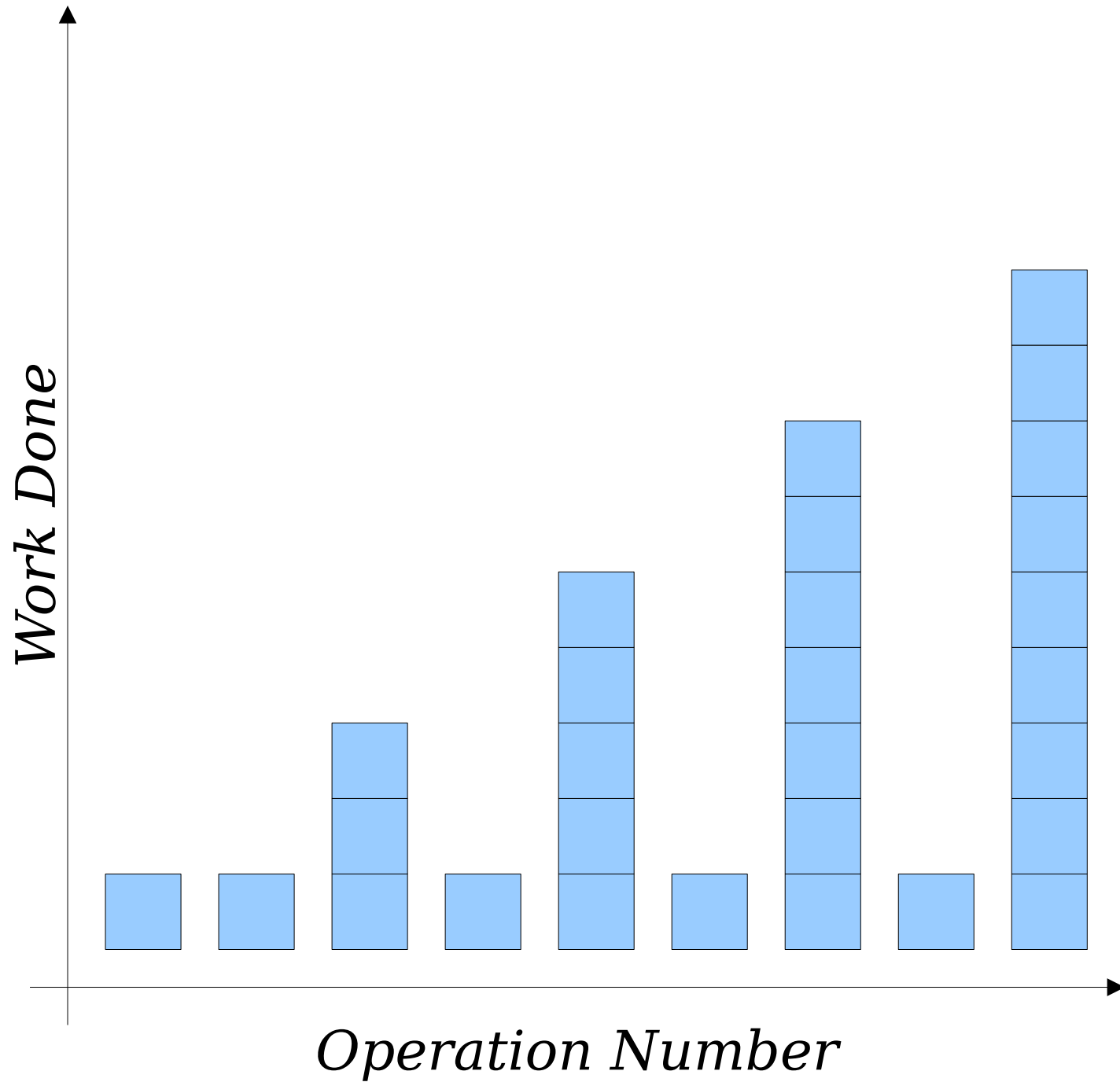




Now, only half the pushes we do will require moving everything to a new array.



Increase array size by *adding one*.



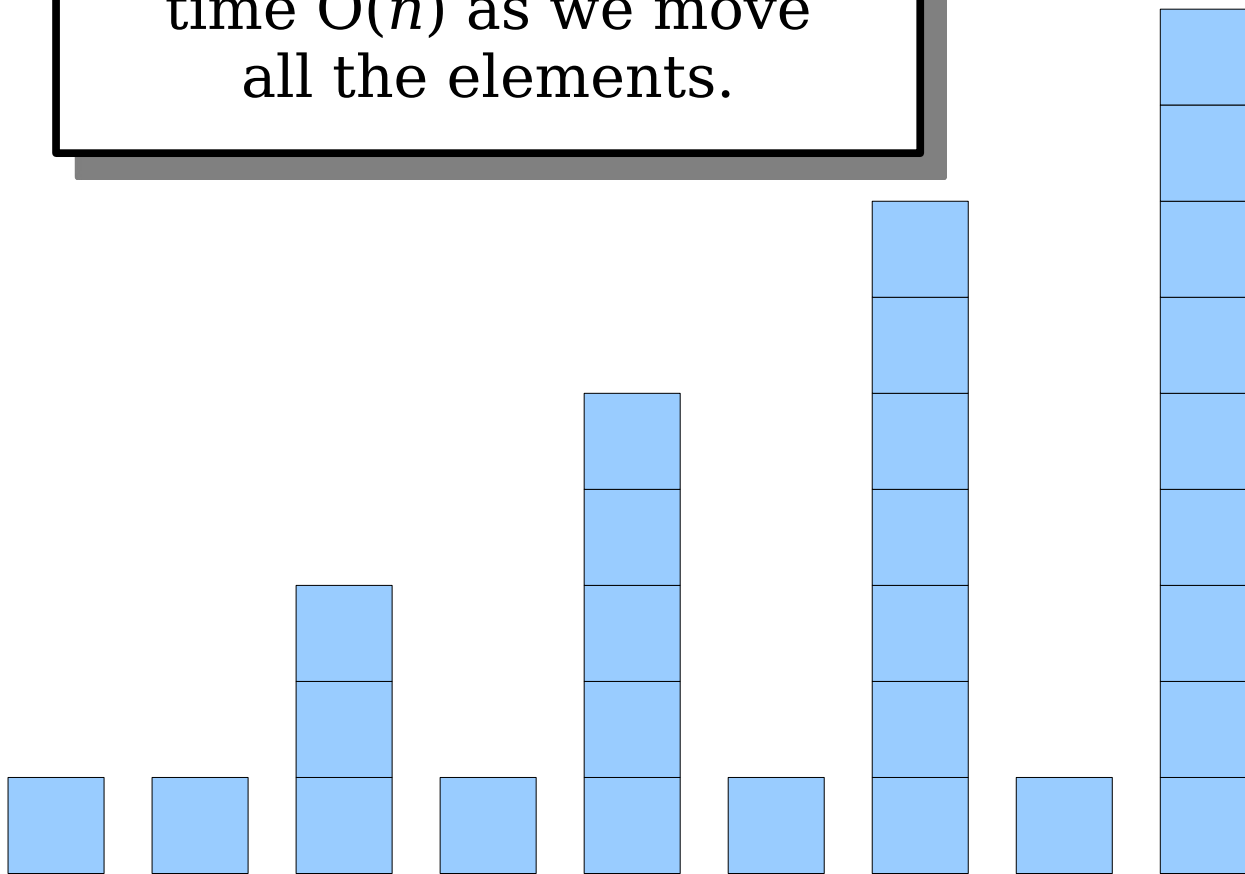
Increase array size by *adding two*.

Half of our pushes take time  $O(1)$  because there's free space left.

Half of our pushes take time  $O(n)$  as we move all the elements.

Increase array size by **adding two**.

*Work Done*

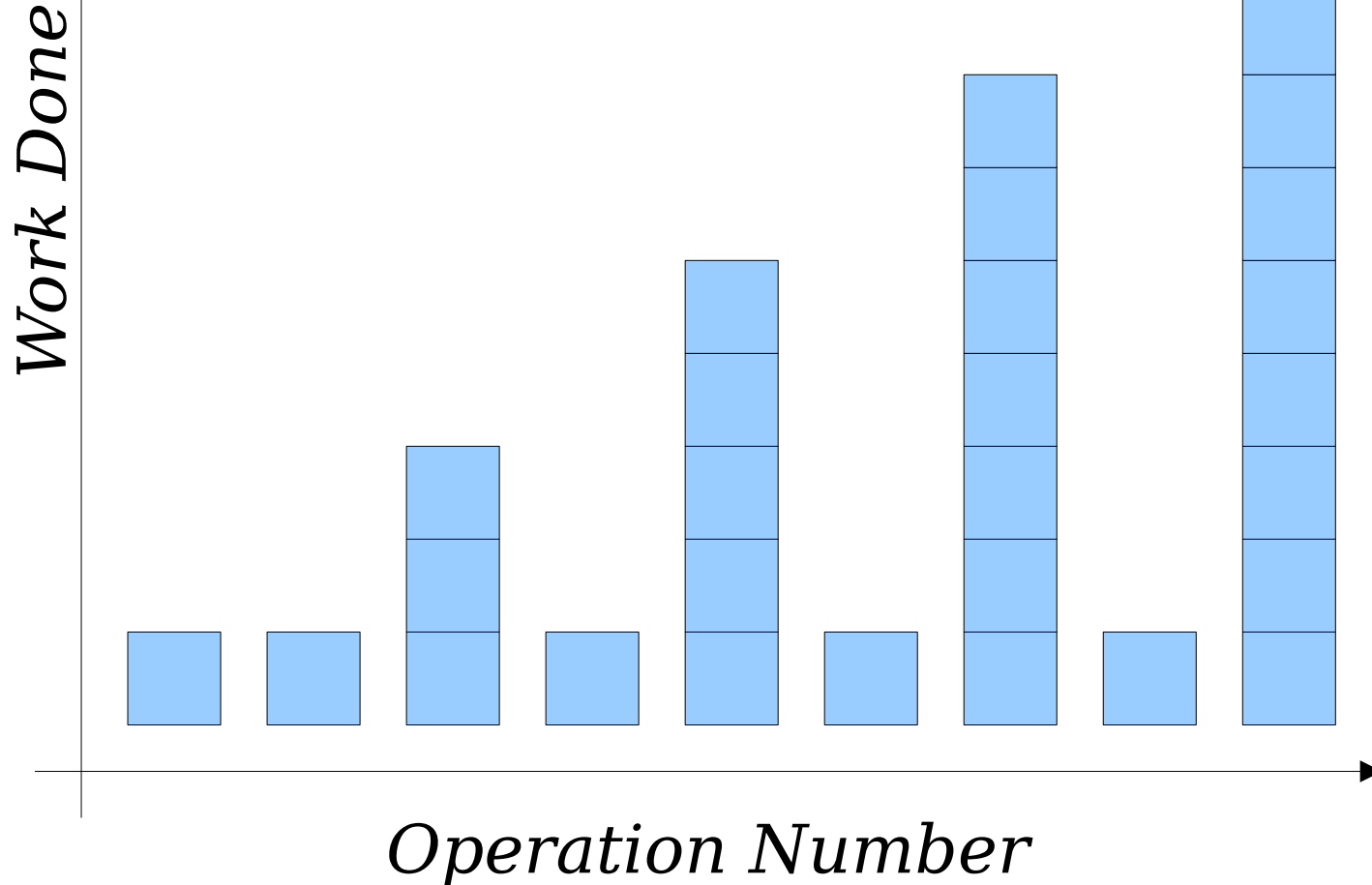


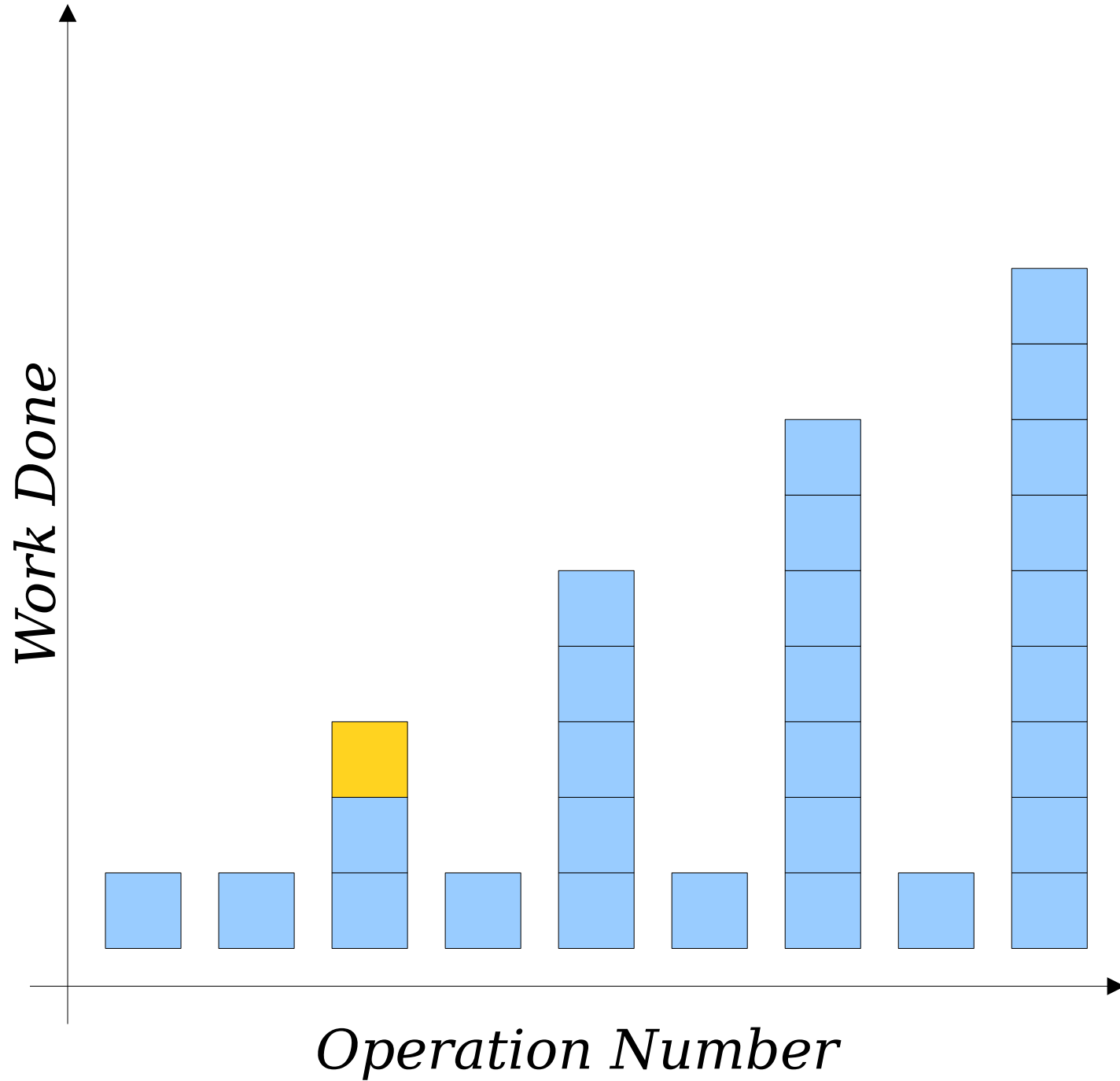
*Operation Number*

What's the average work done with each push?

To find out, let's see how much total work was done.

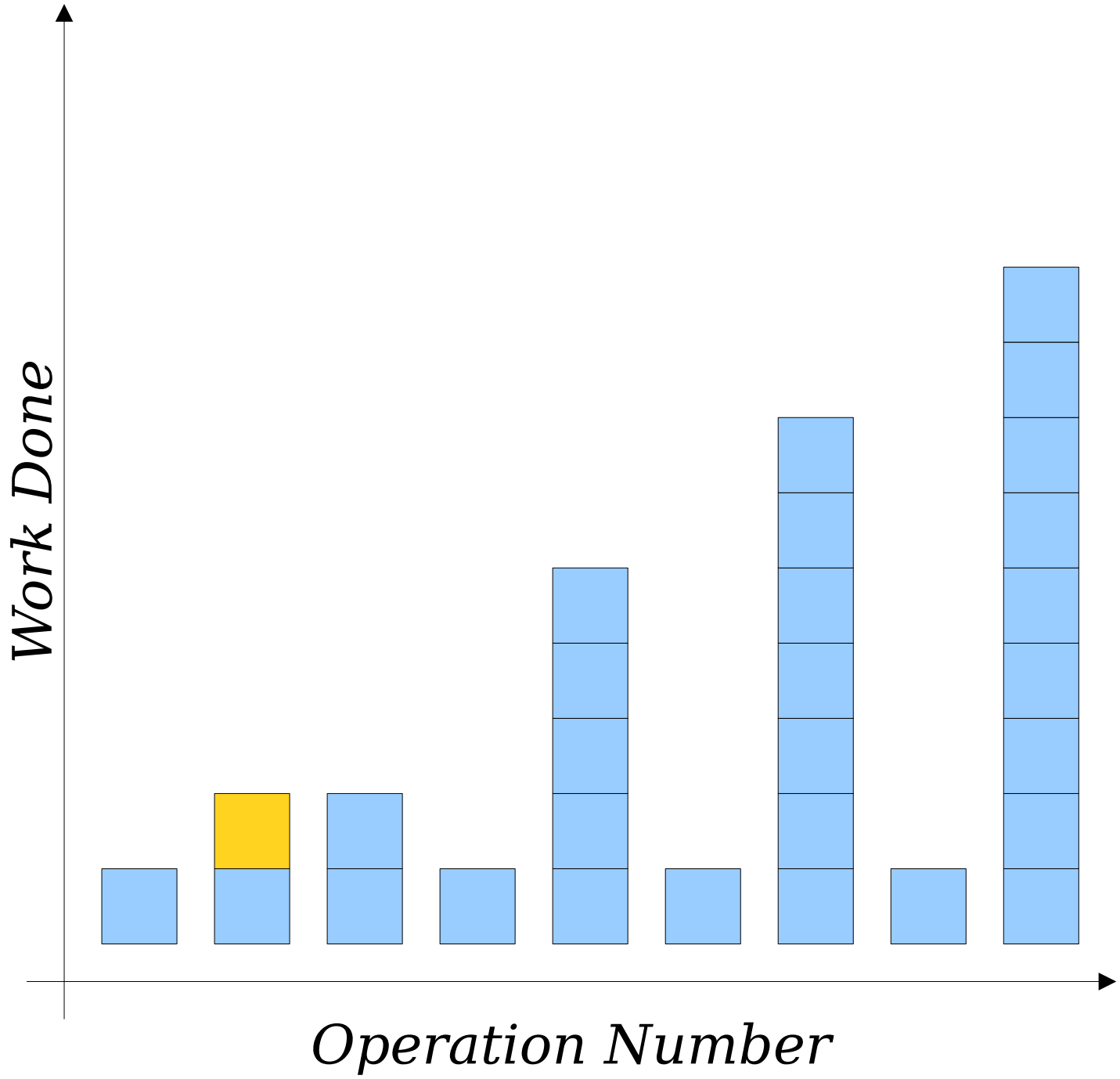
Increase array size by **adding two**.



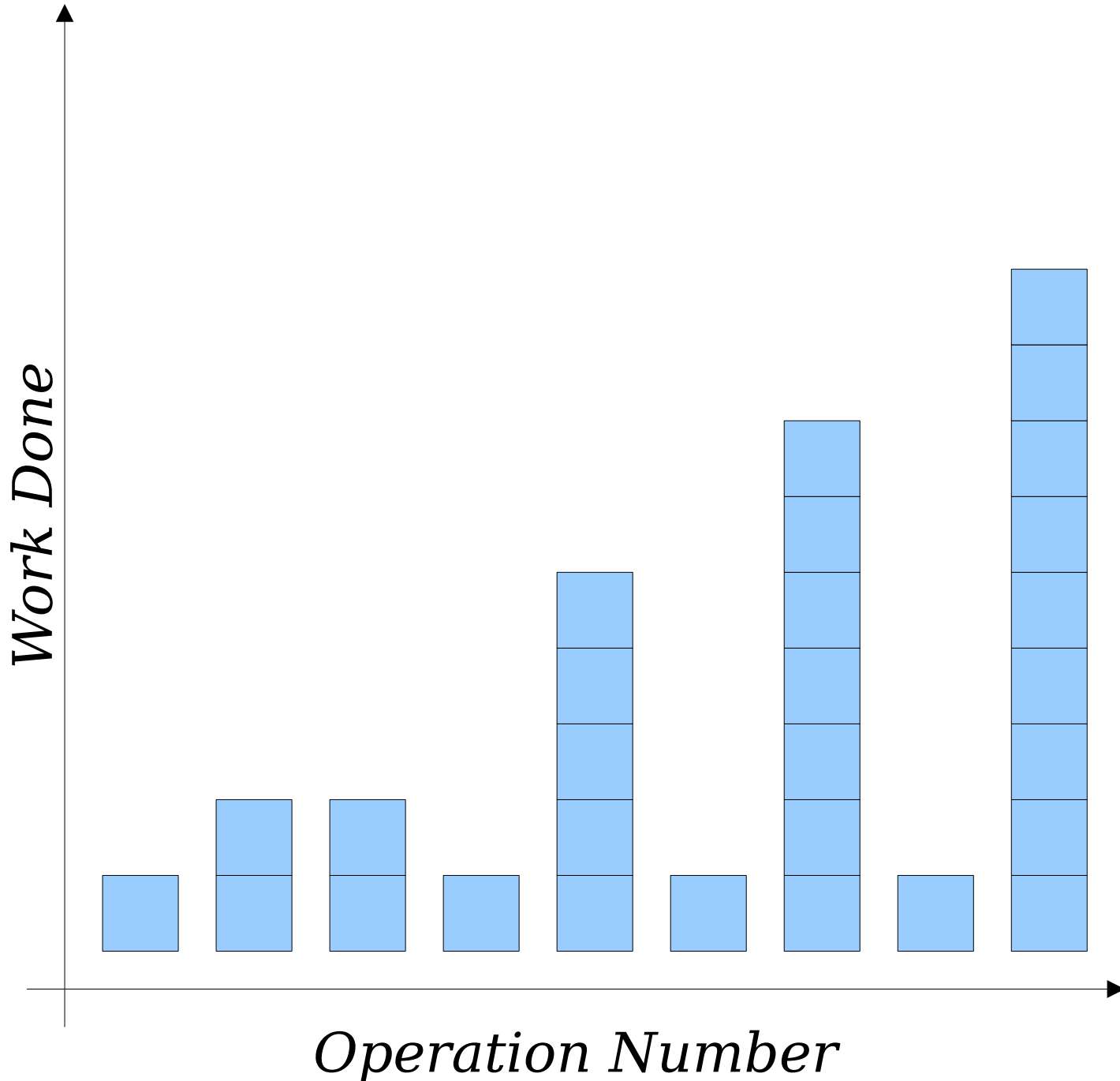


Increase array size by *adding two*.

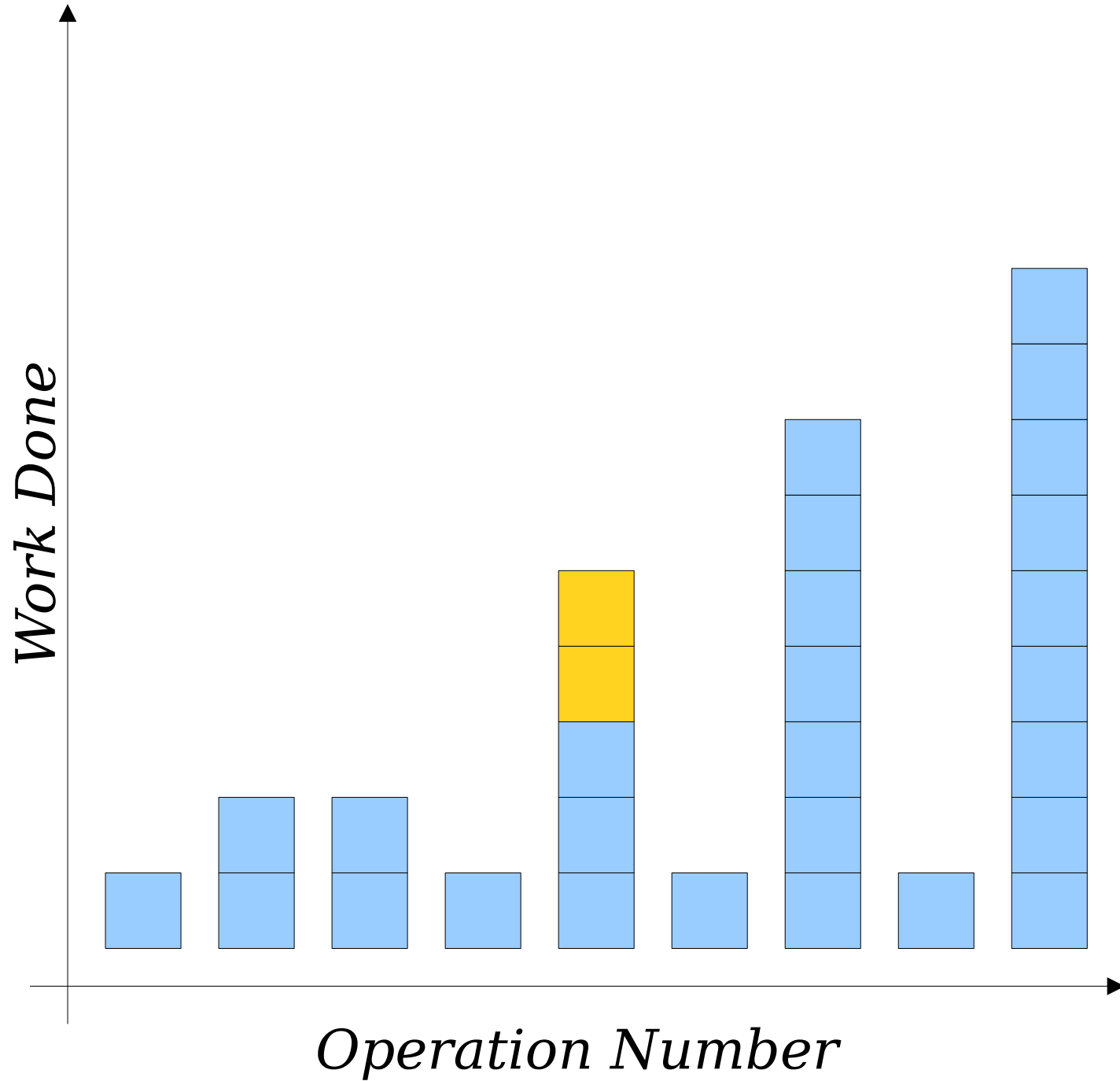
Increase array size by *adding two*.



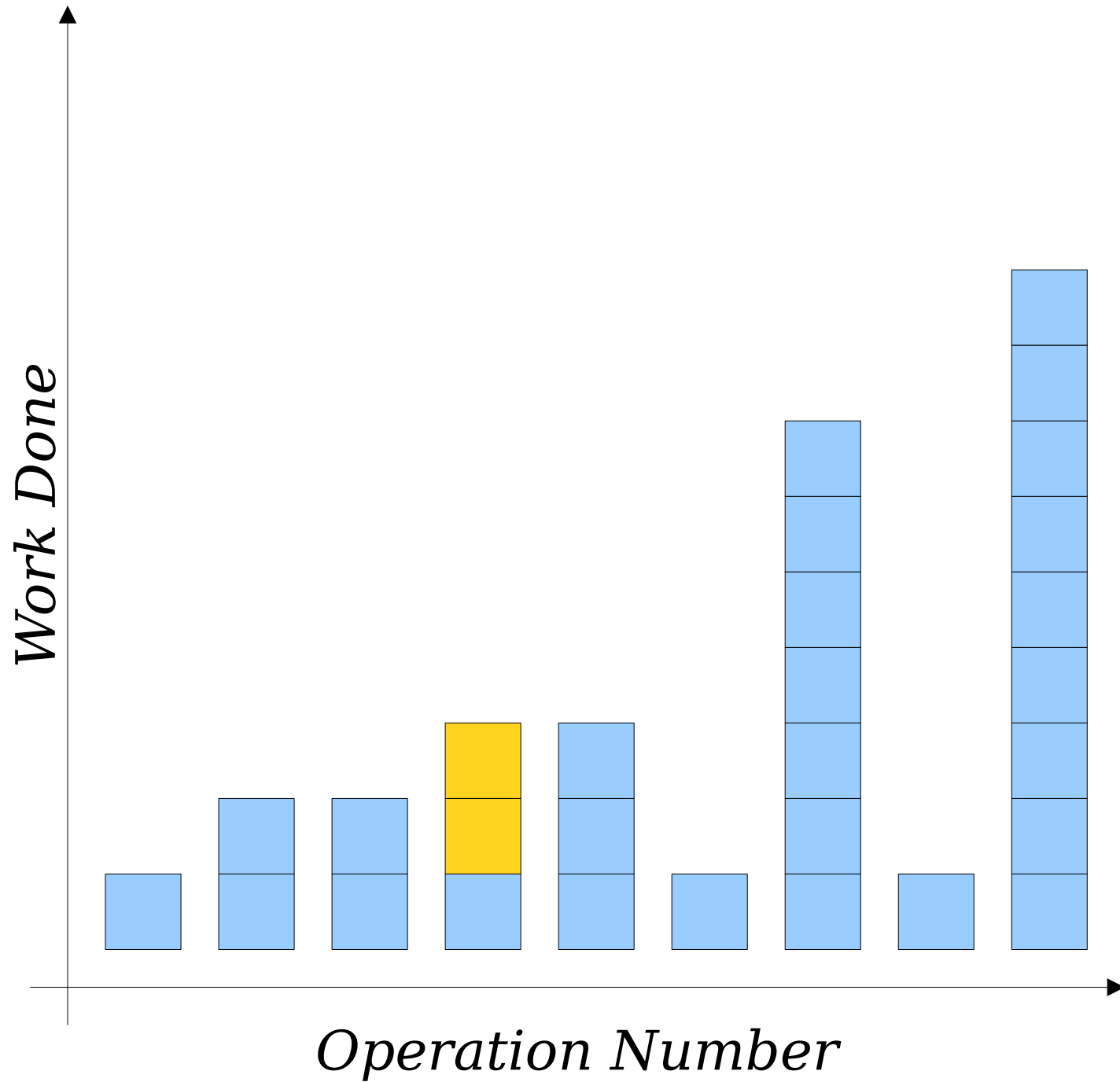
Increase array size by *adding two*.



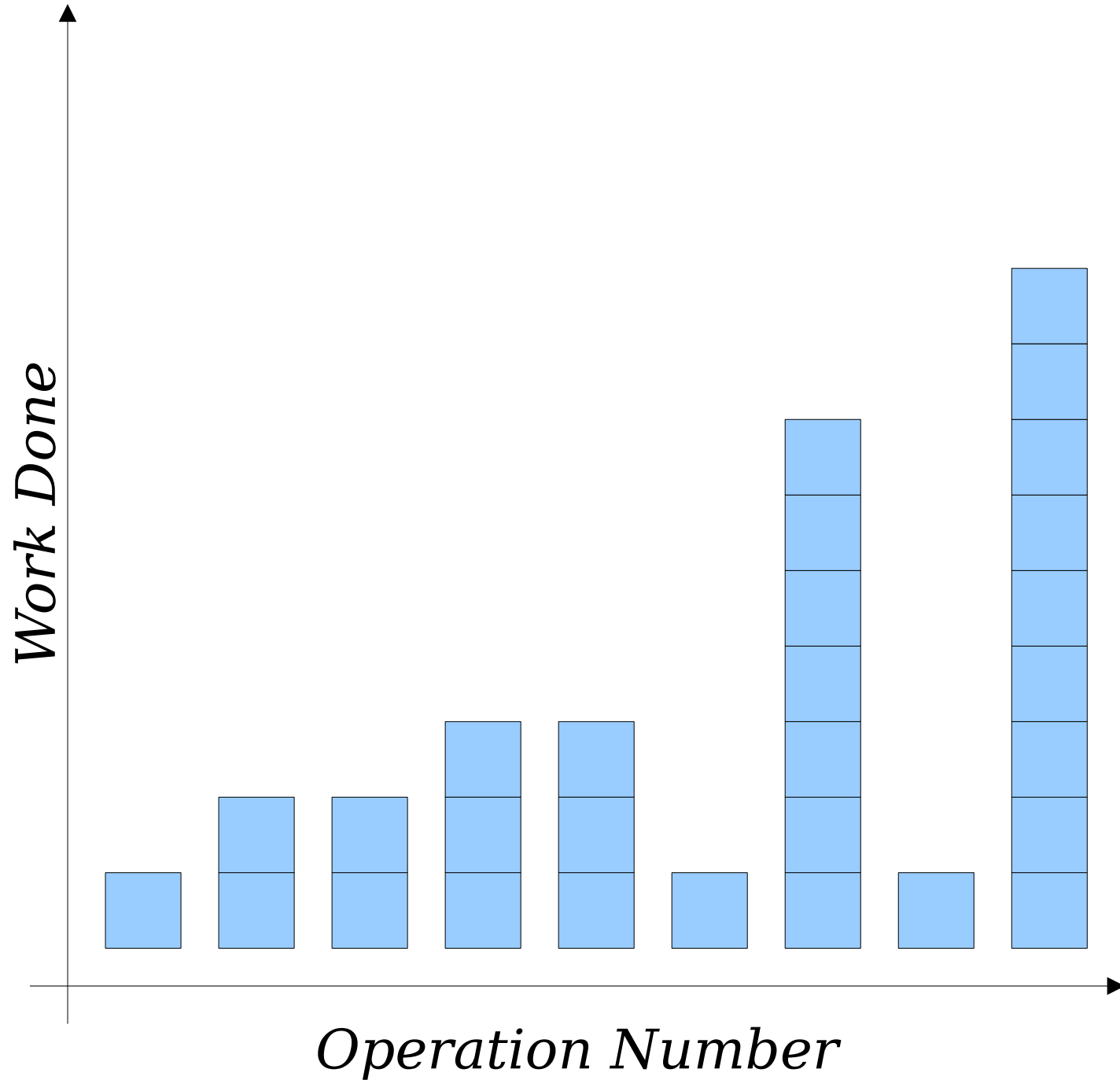




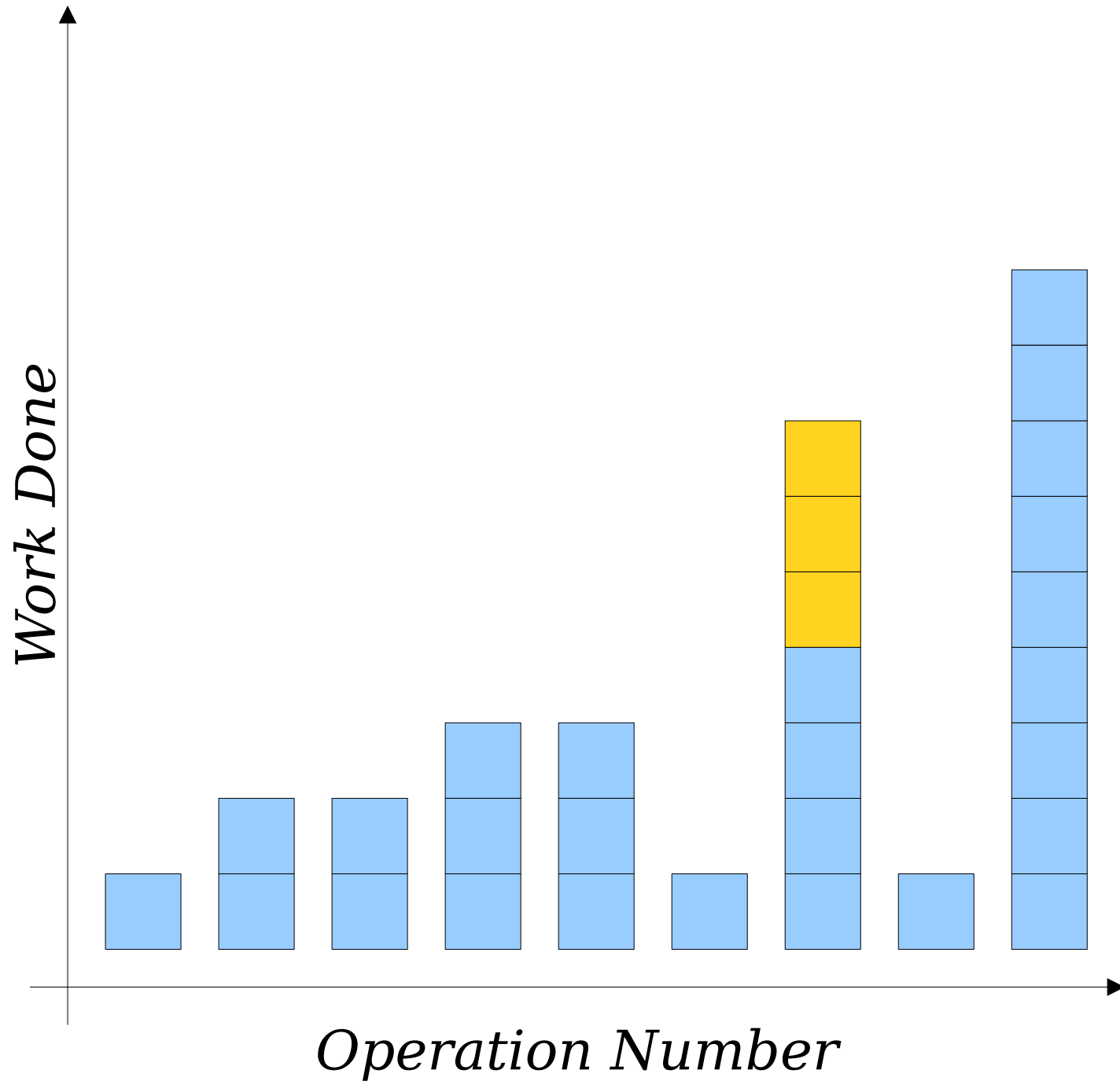
Increase array size by **adding two**.



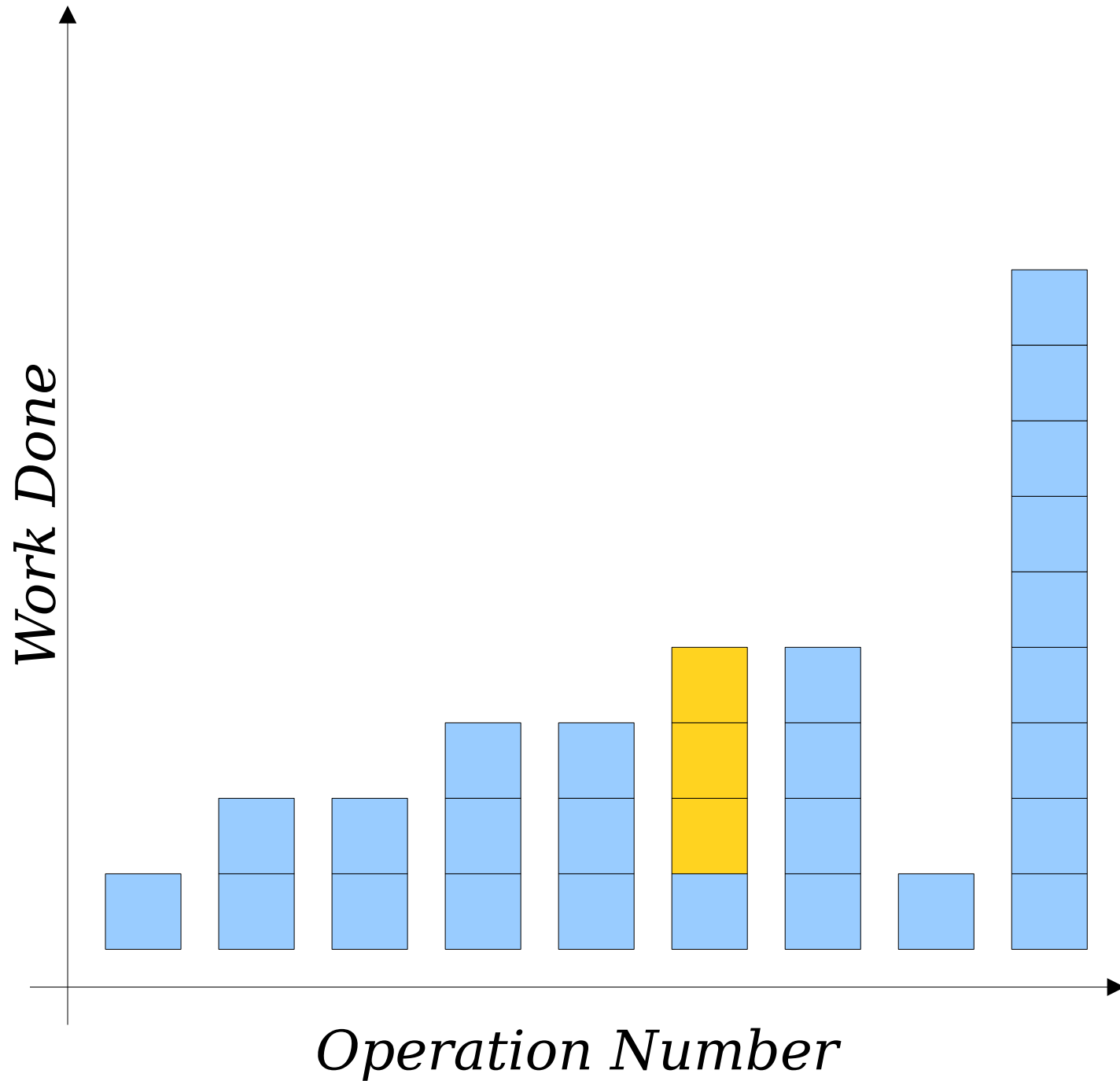
Increase array size by *adding two*.



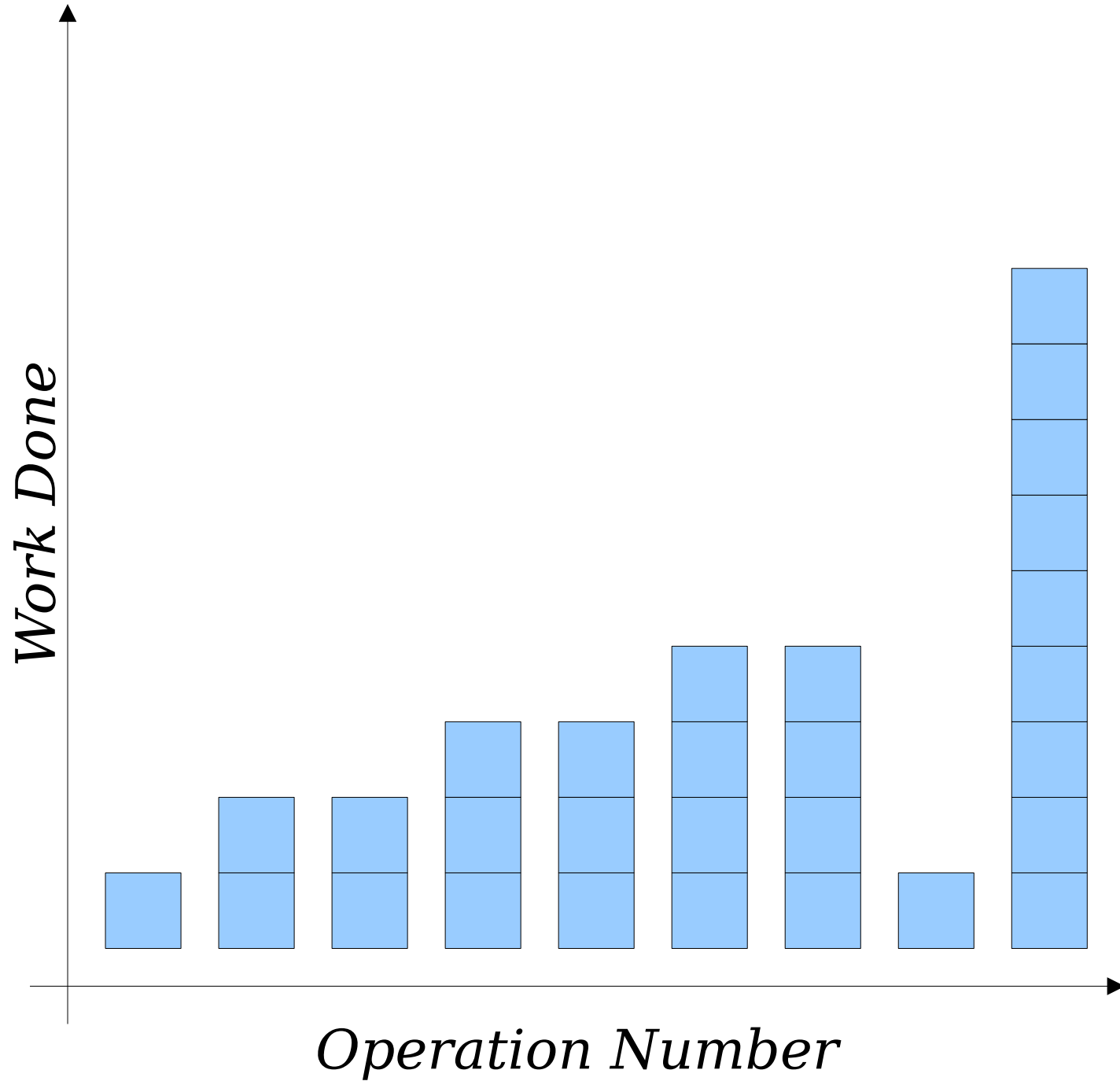
Increase array size by *adding two*.



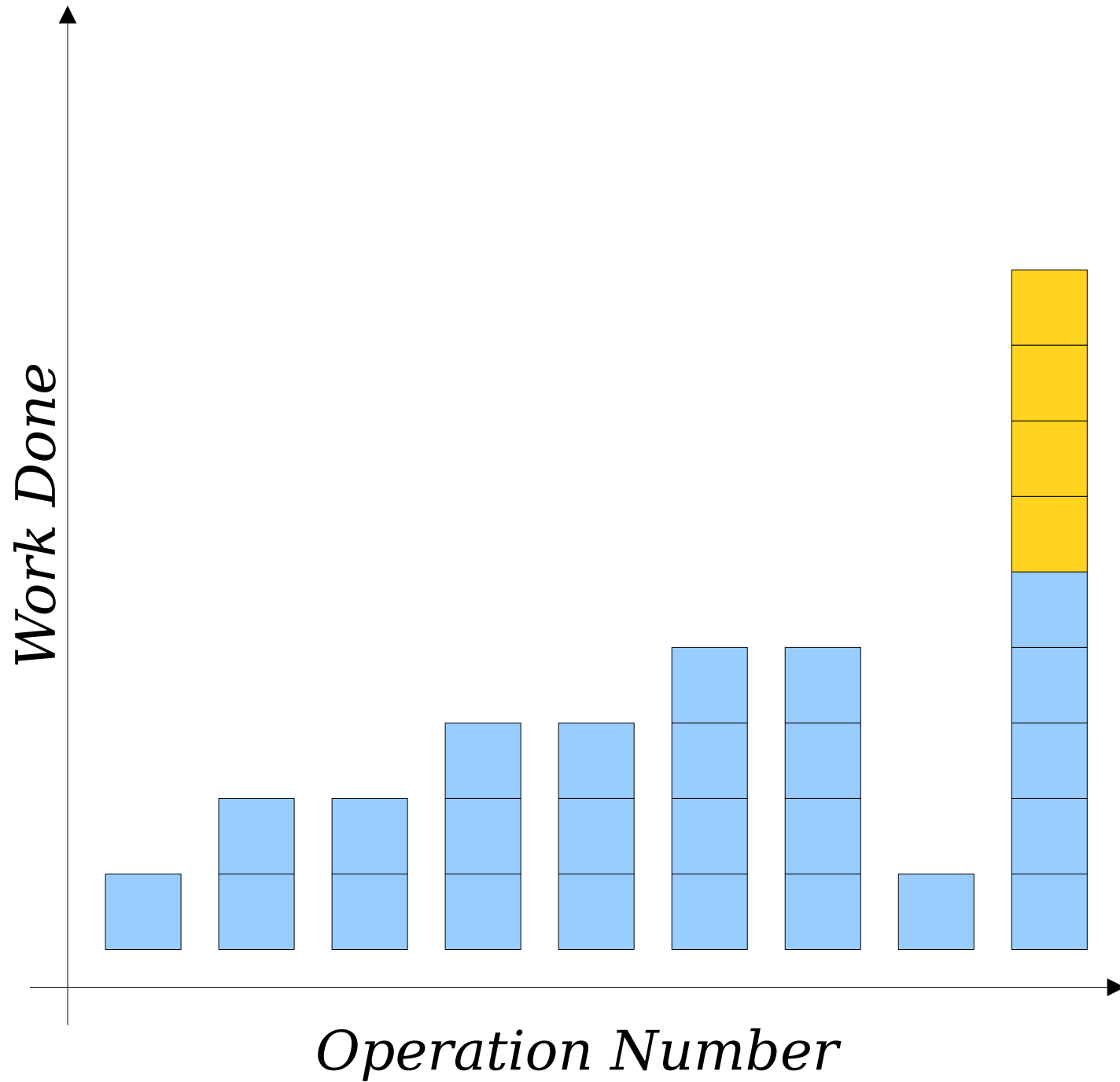
Increase array size by *adding two*.



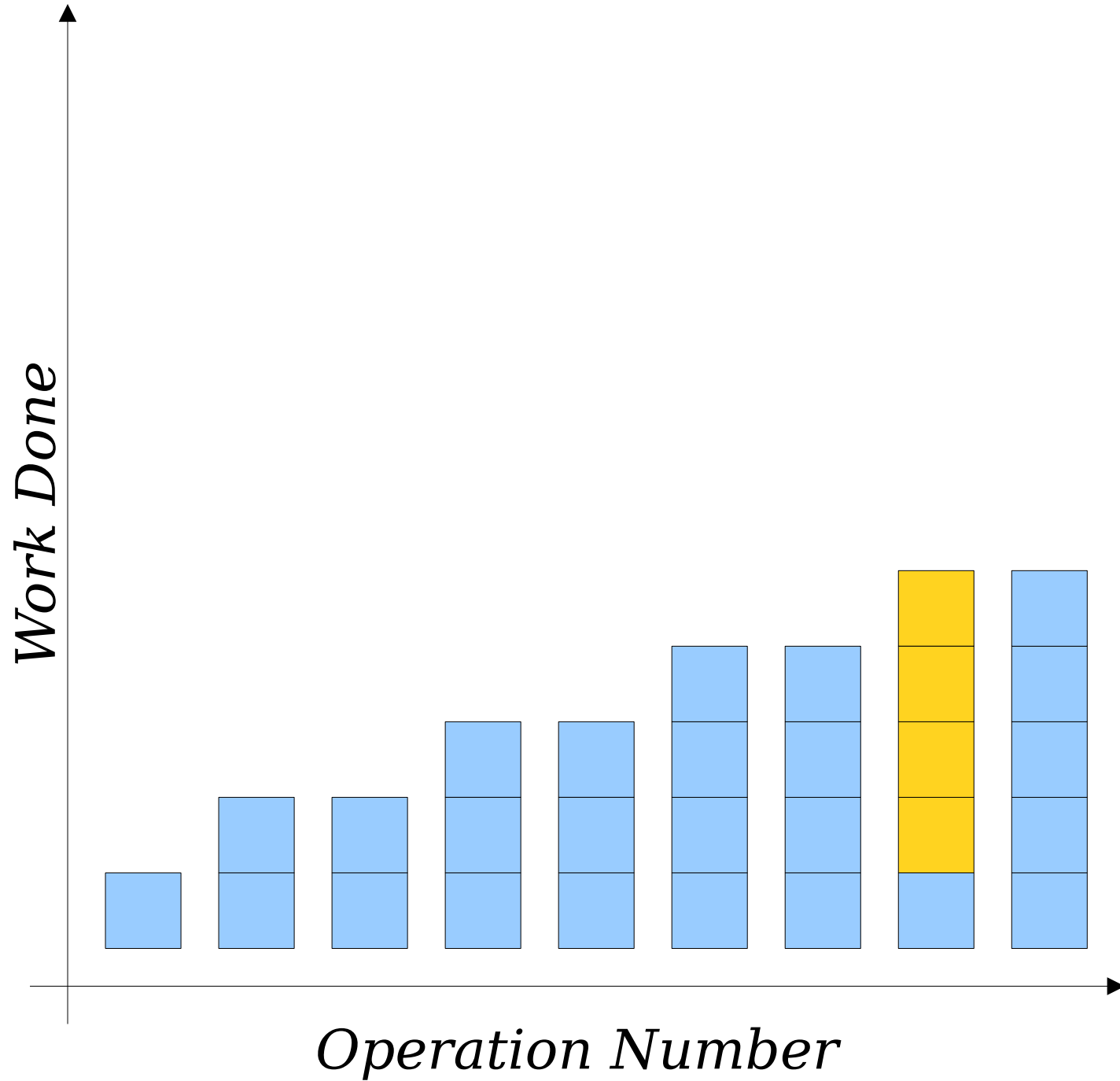
Increase array size by **adding two**.



Increase array size by *adding two*.



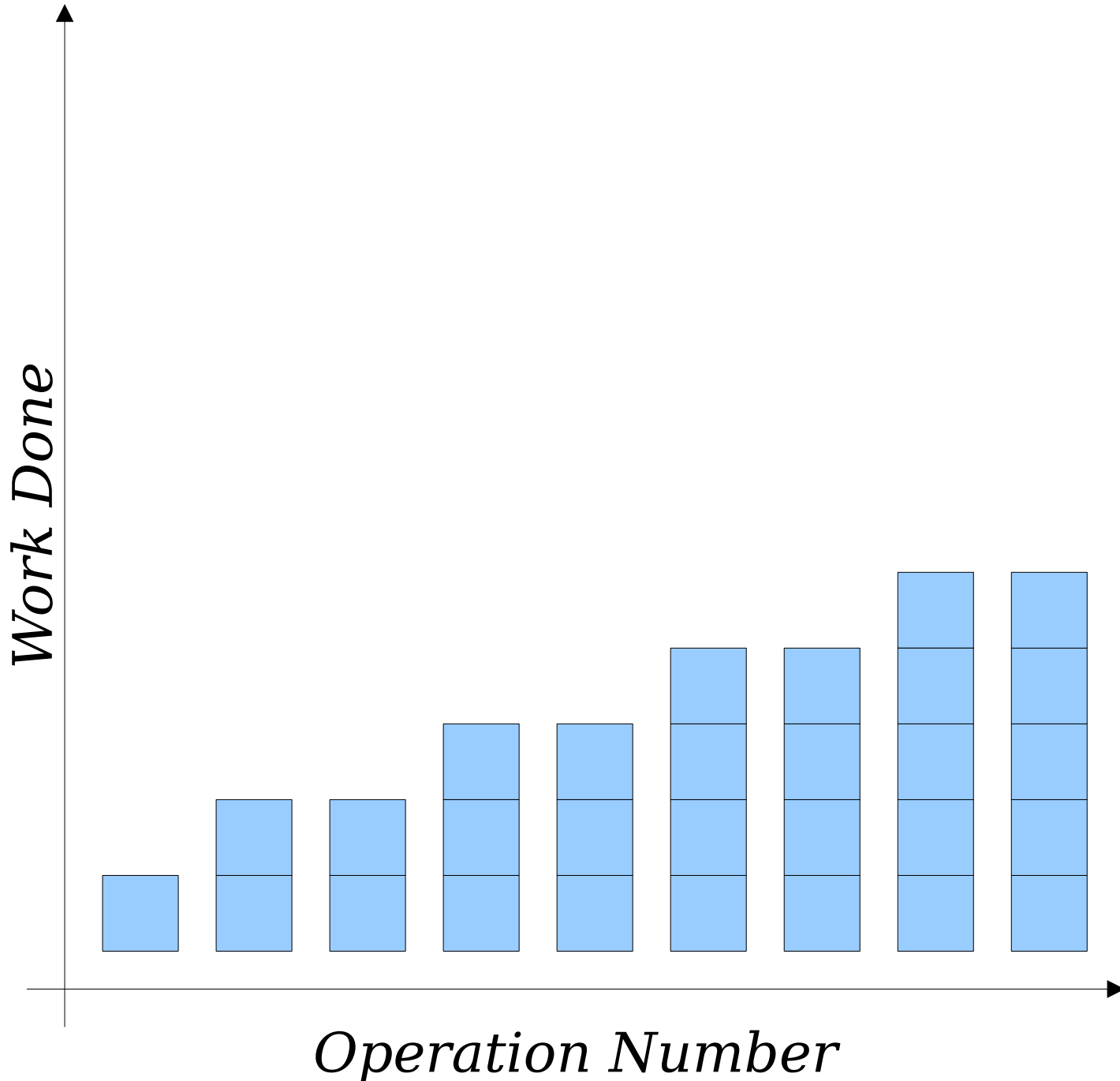
Increase array size by *adding two*.

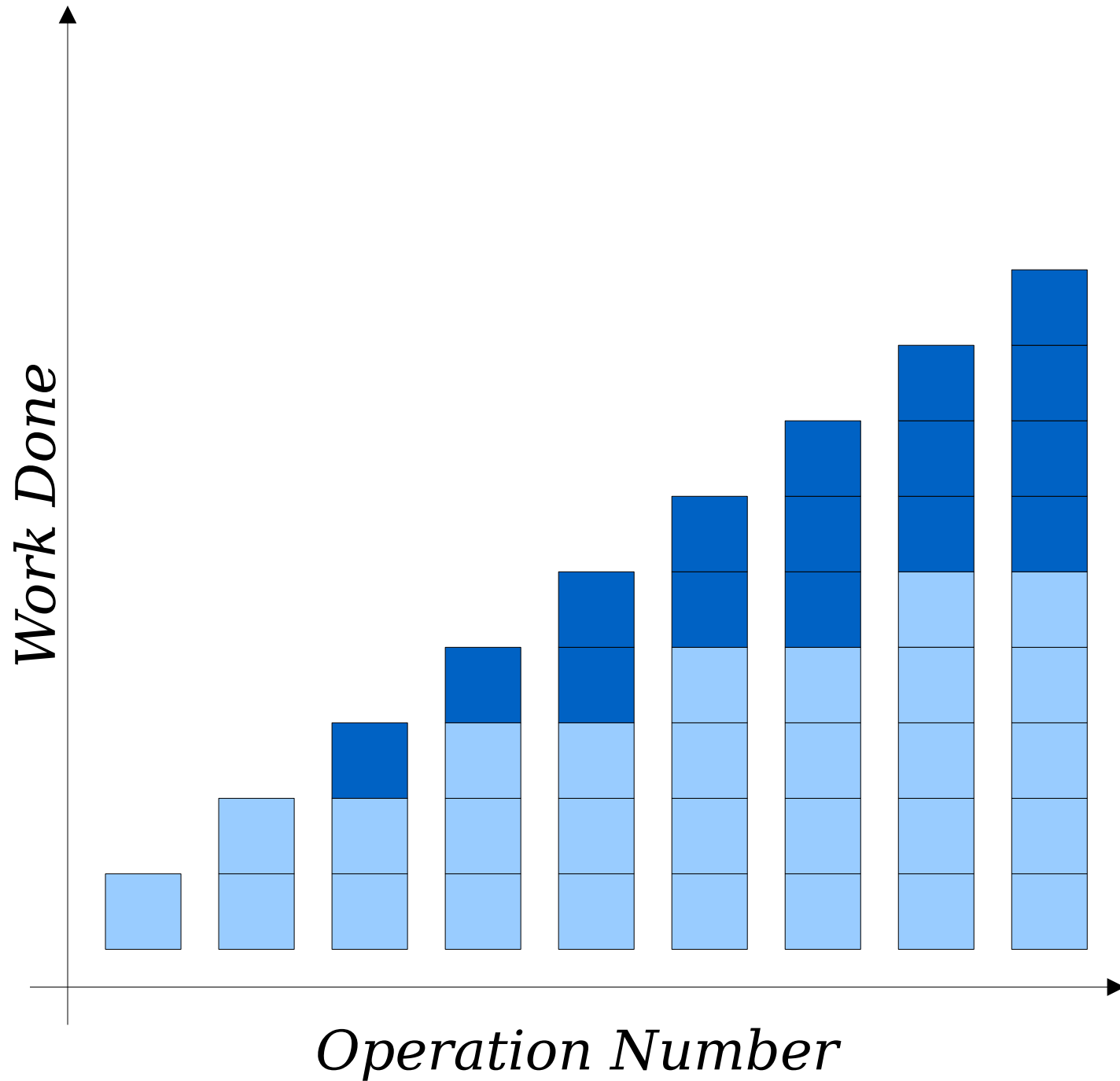


Increase array size by *adding two*.



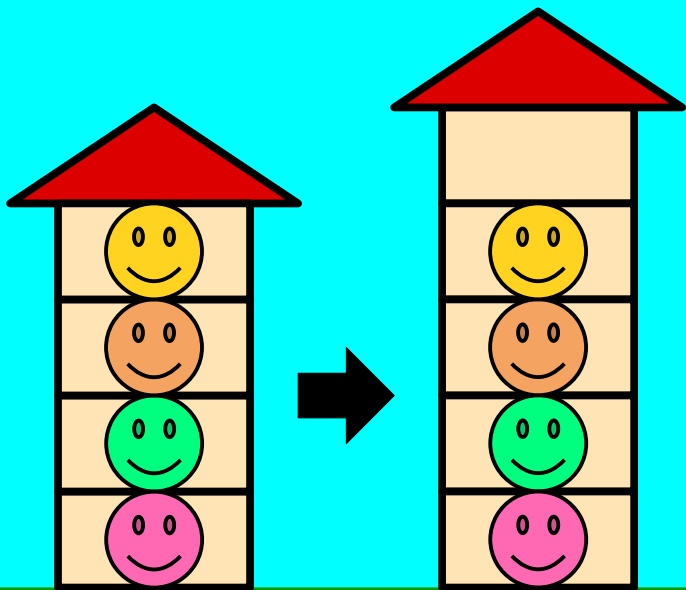
Increase array size by *adding two*.

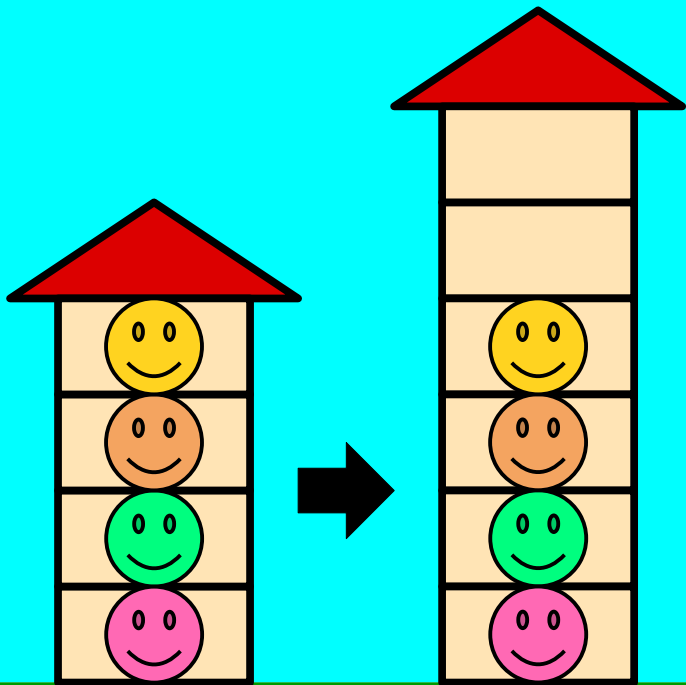


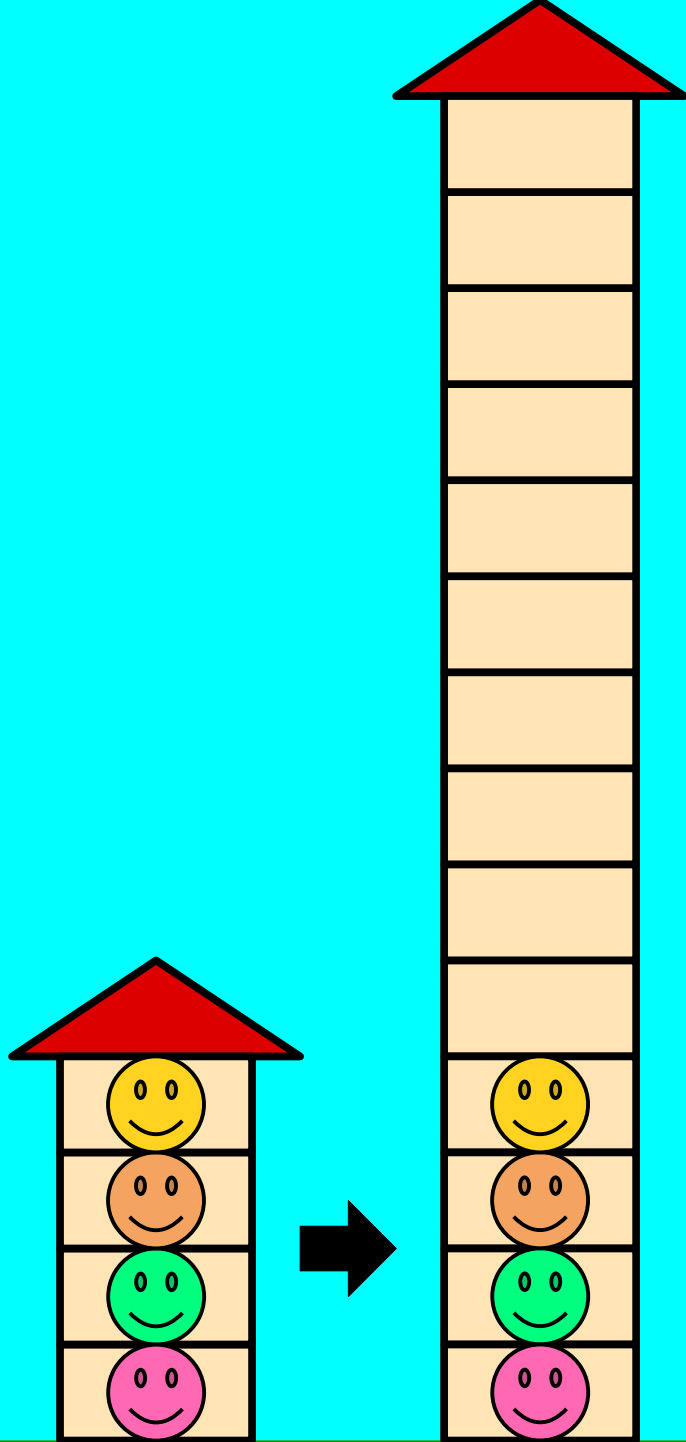


Increase array size by *adding two*.

This roughly halves the work done.

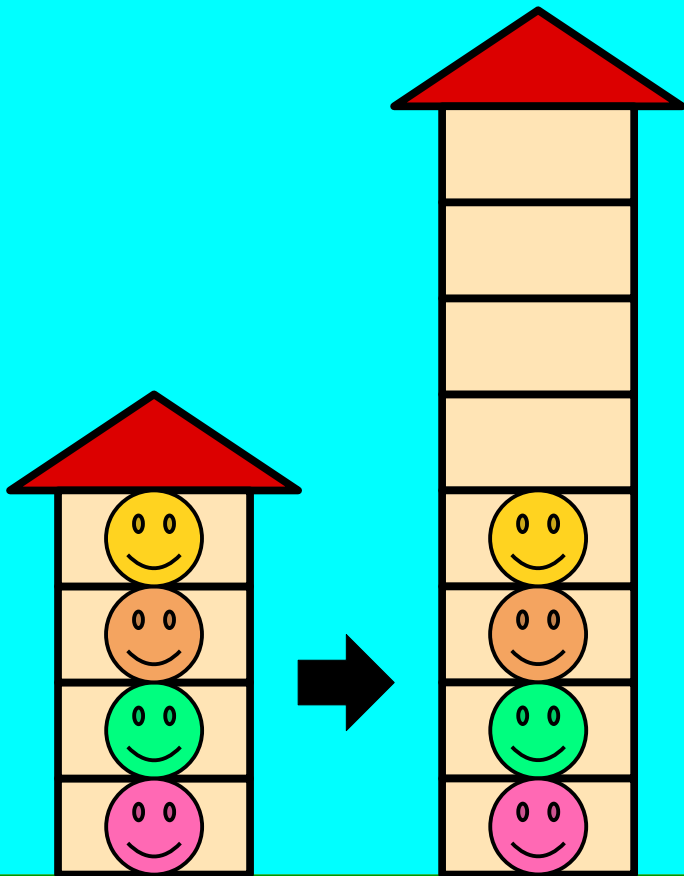






If we make the new array too big, we're might not make use of all the new space.

What's a good compromise?



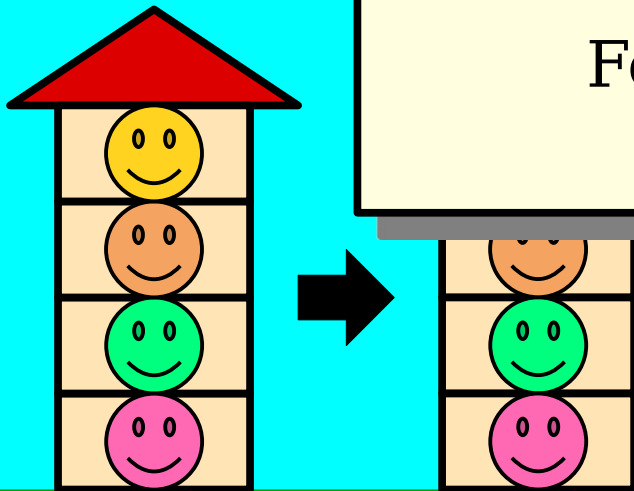
**Idea:** Make the new array twice as big as the old one.

This gives us a lot of free space, and we never use more than twice the space we need.

What is the big-O cost of a push?  
What is the big-O cost of  $n$  pushes?

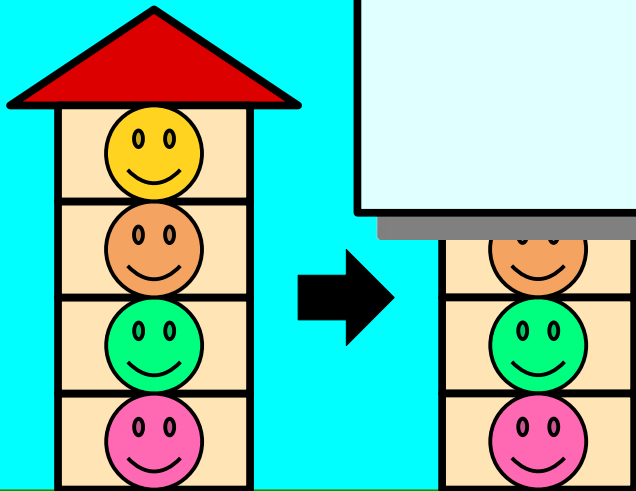
Take thirty seconds to think this over.  
Look over the runtime plot and the  
code for the push operation.

Formulate a hypothesis!



What is the big-O cost of a push?  
What is the big-O cost of  $n$  pushes?

Discuss with your neighbor!



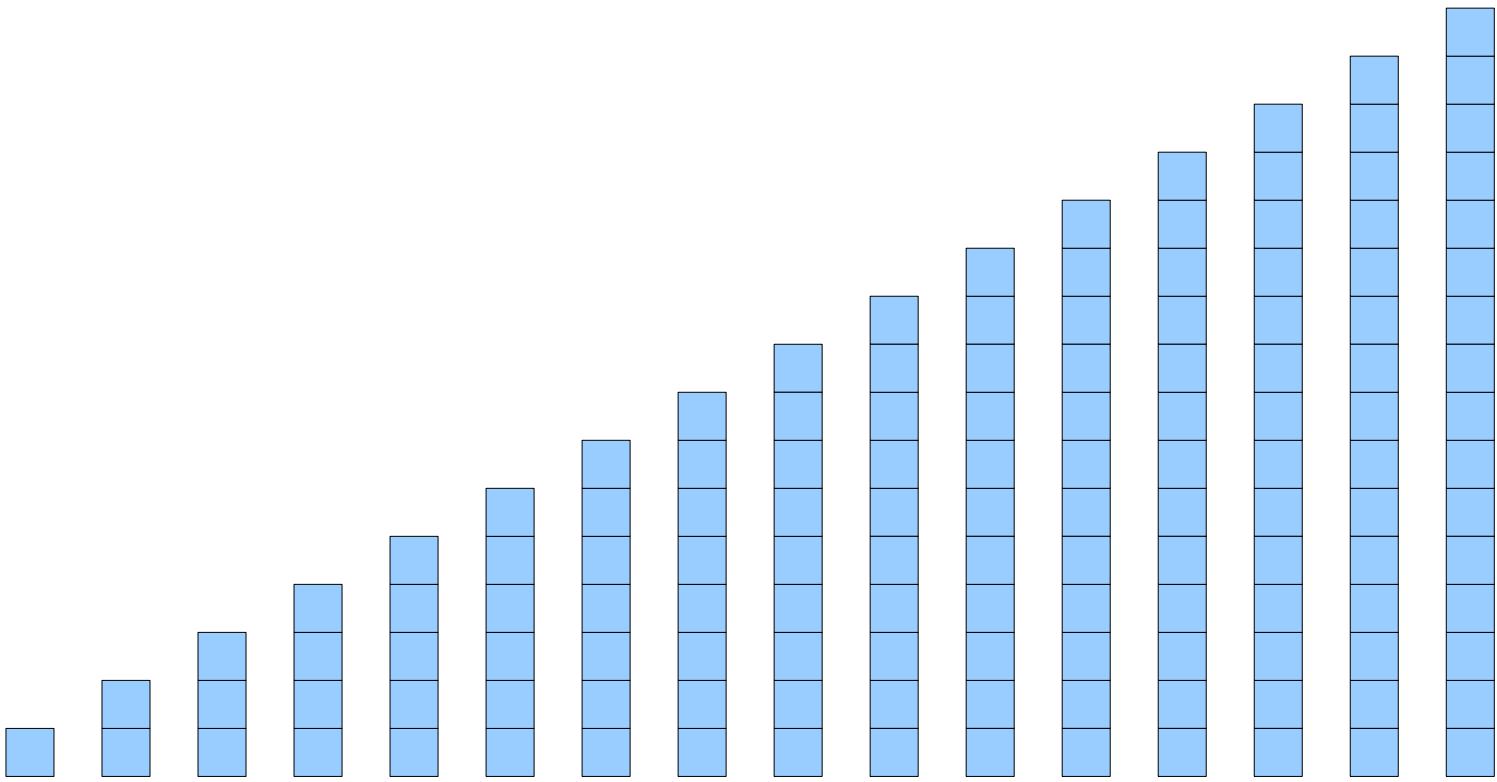


How do we analyze this?

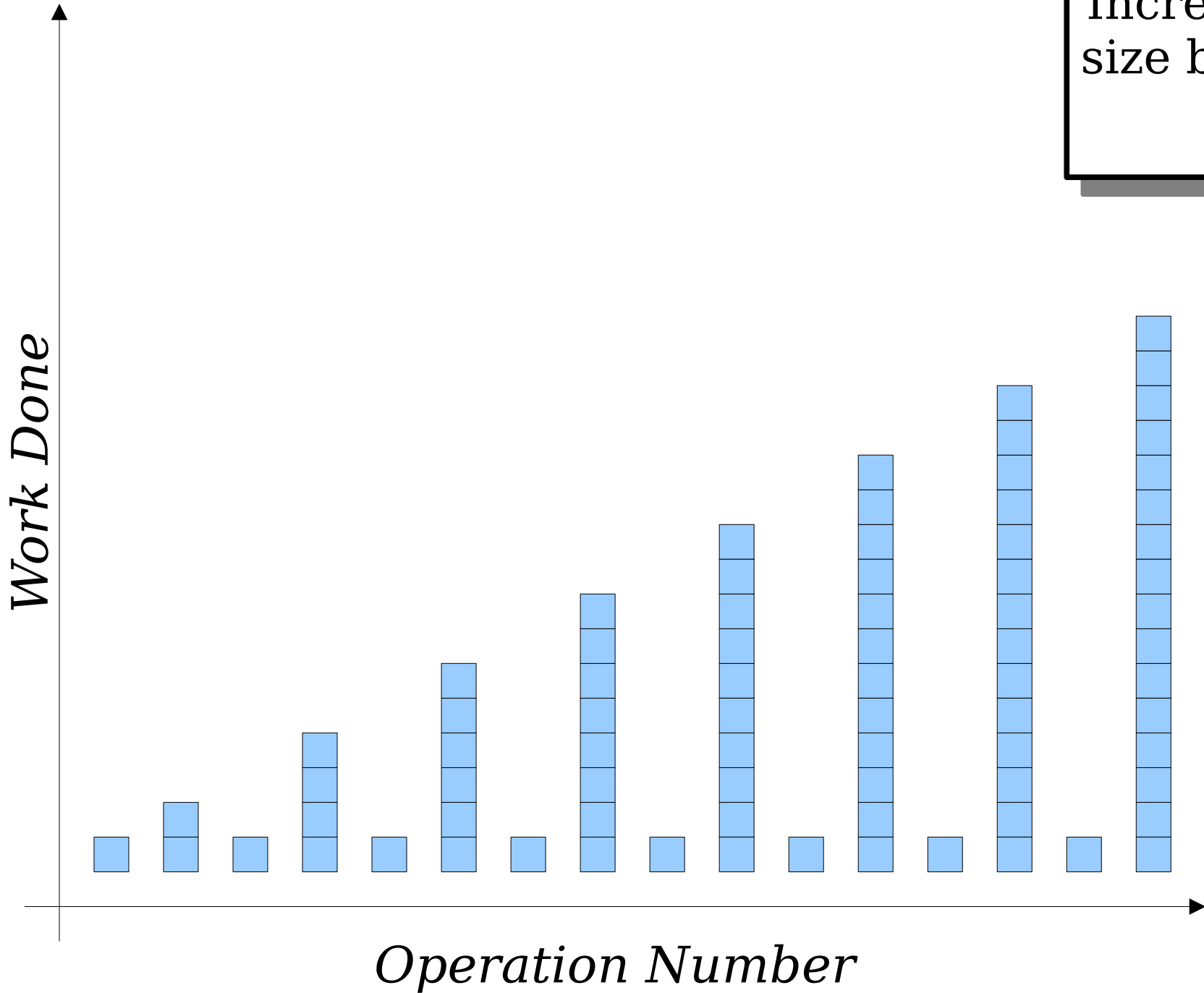
*Work Done*

*Operation Number*

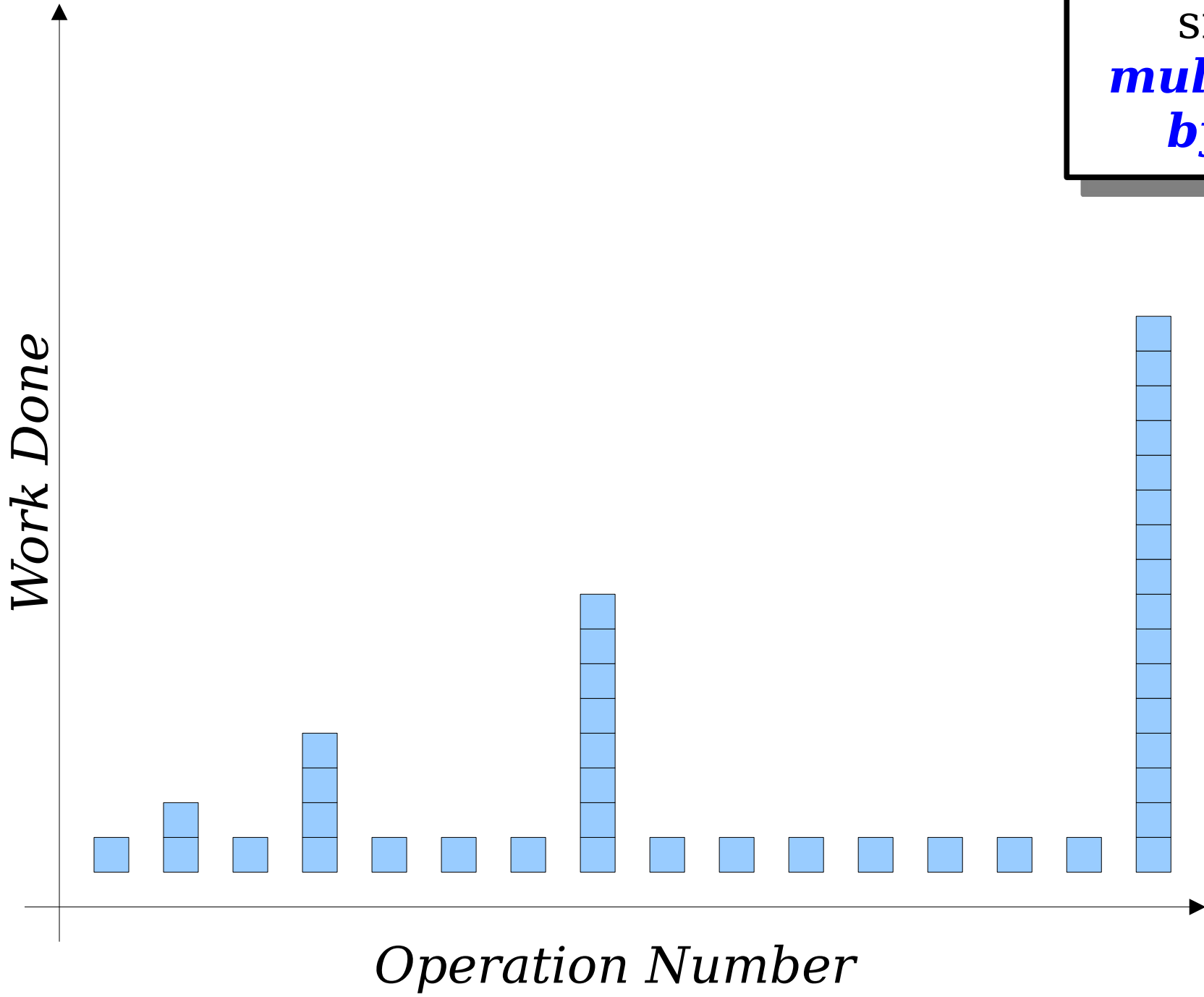
Increase array size by ***adding one.***

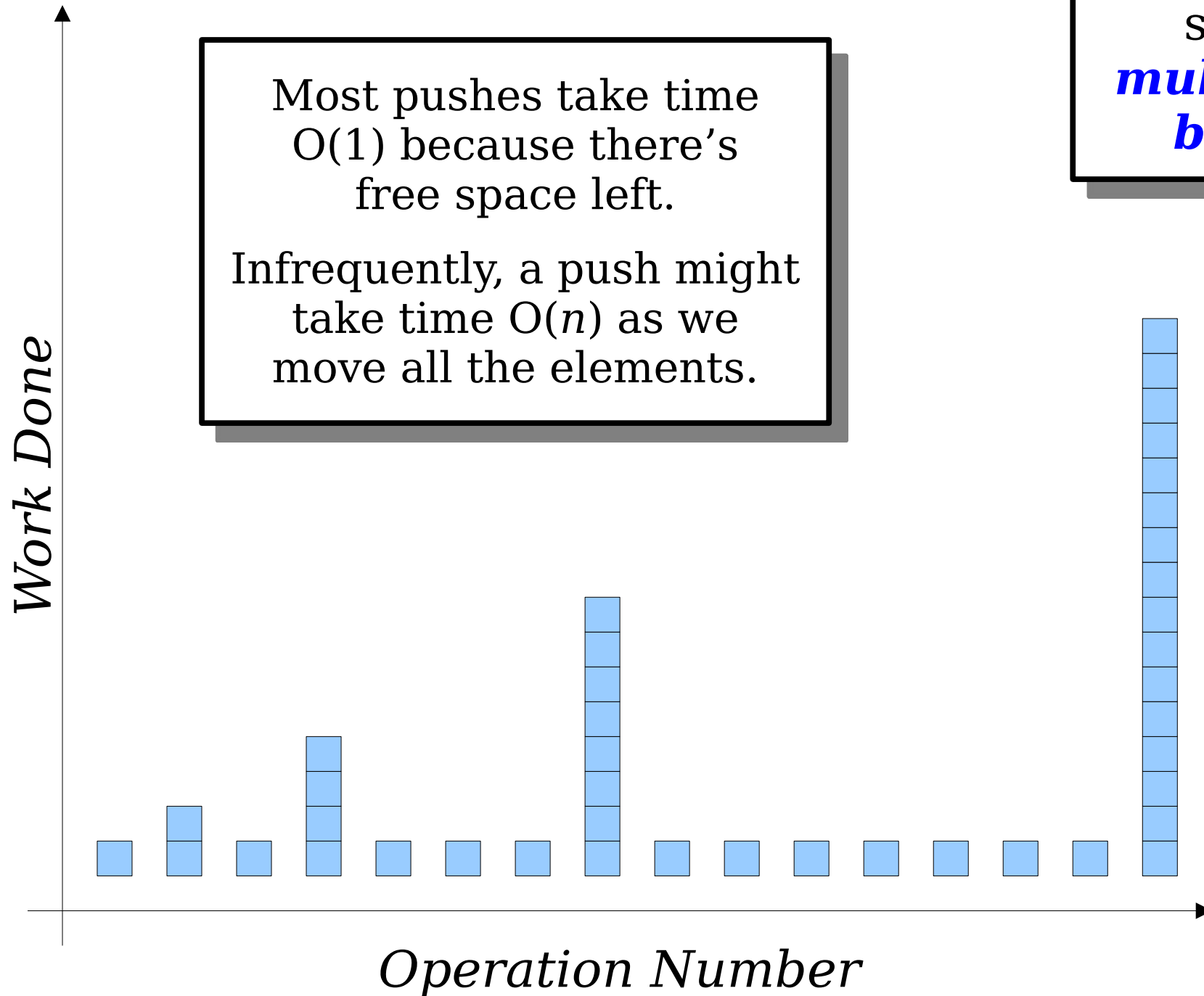


Increase array size by *adding two*.



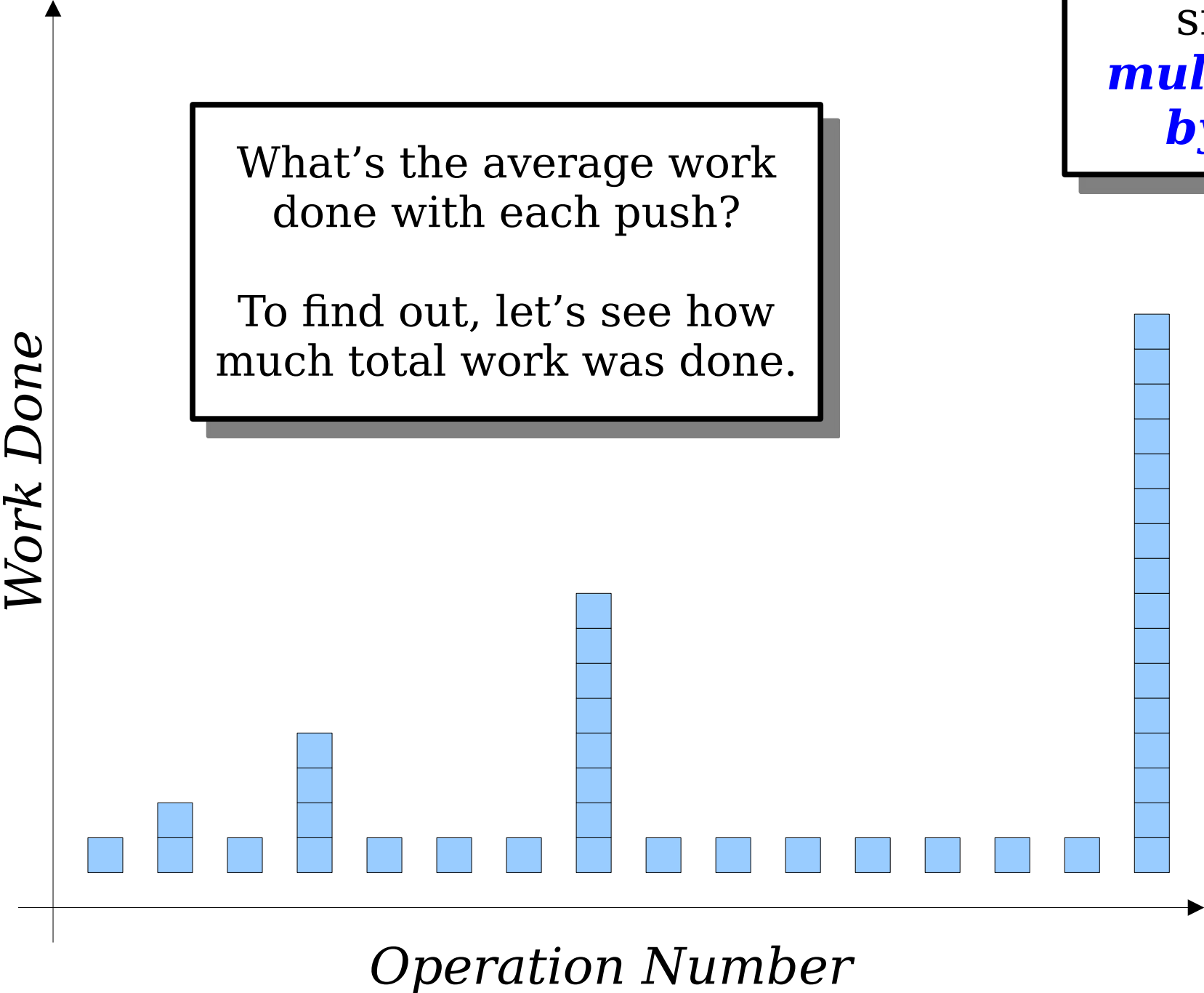
Increase array size by  
size by  
***multiplying  
by two.***



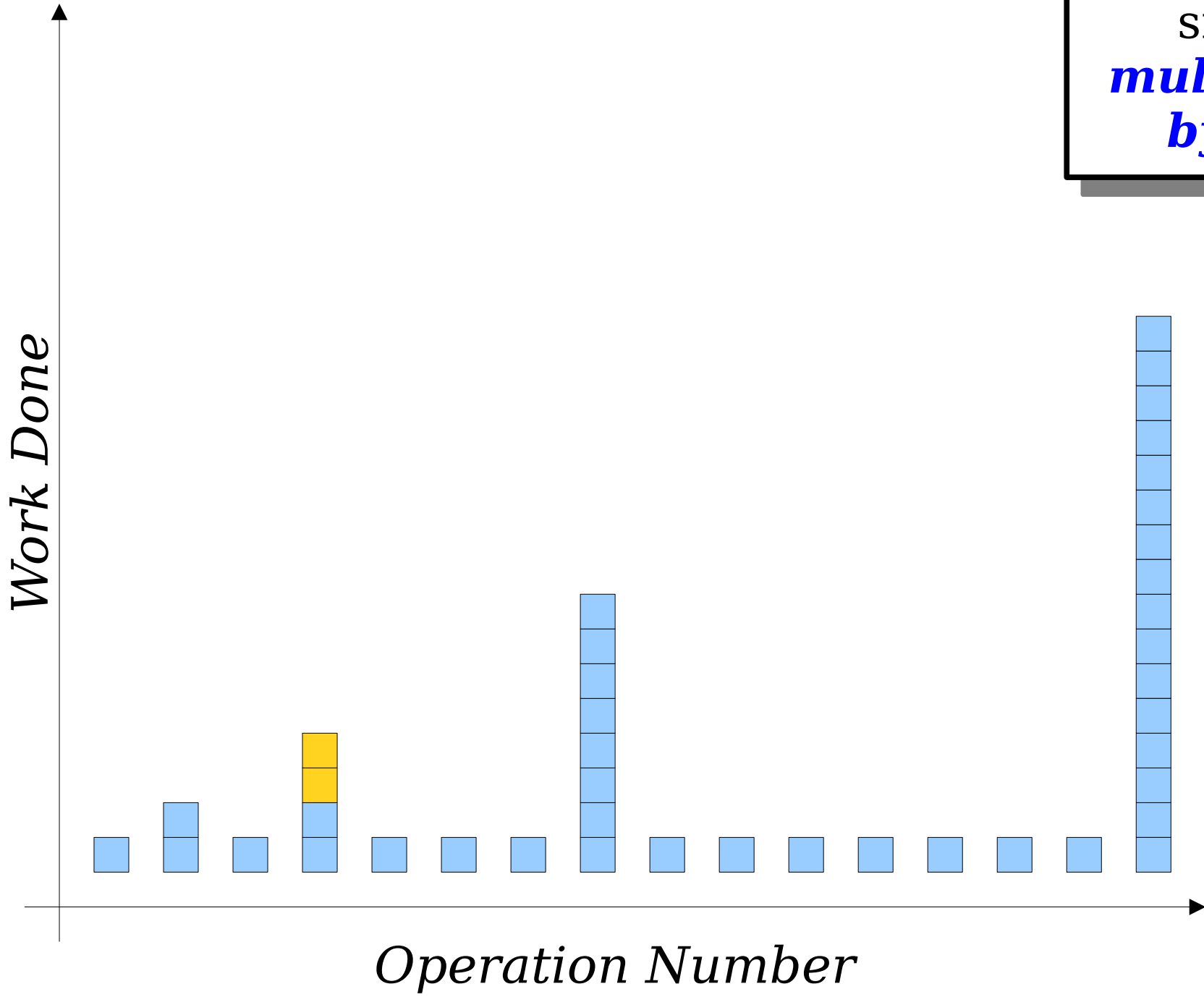


Increase array size by *multiplying by two*.

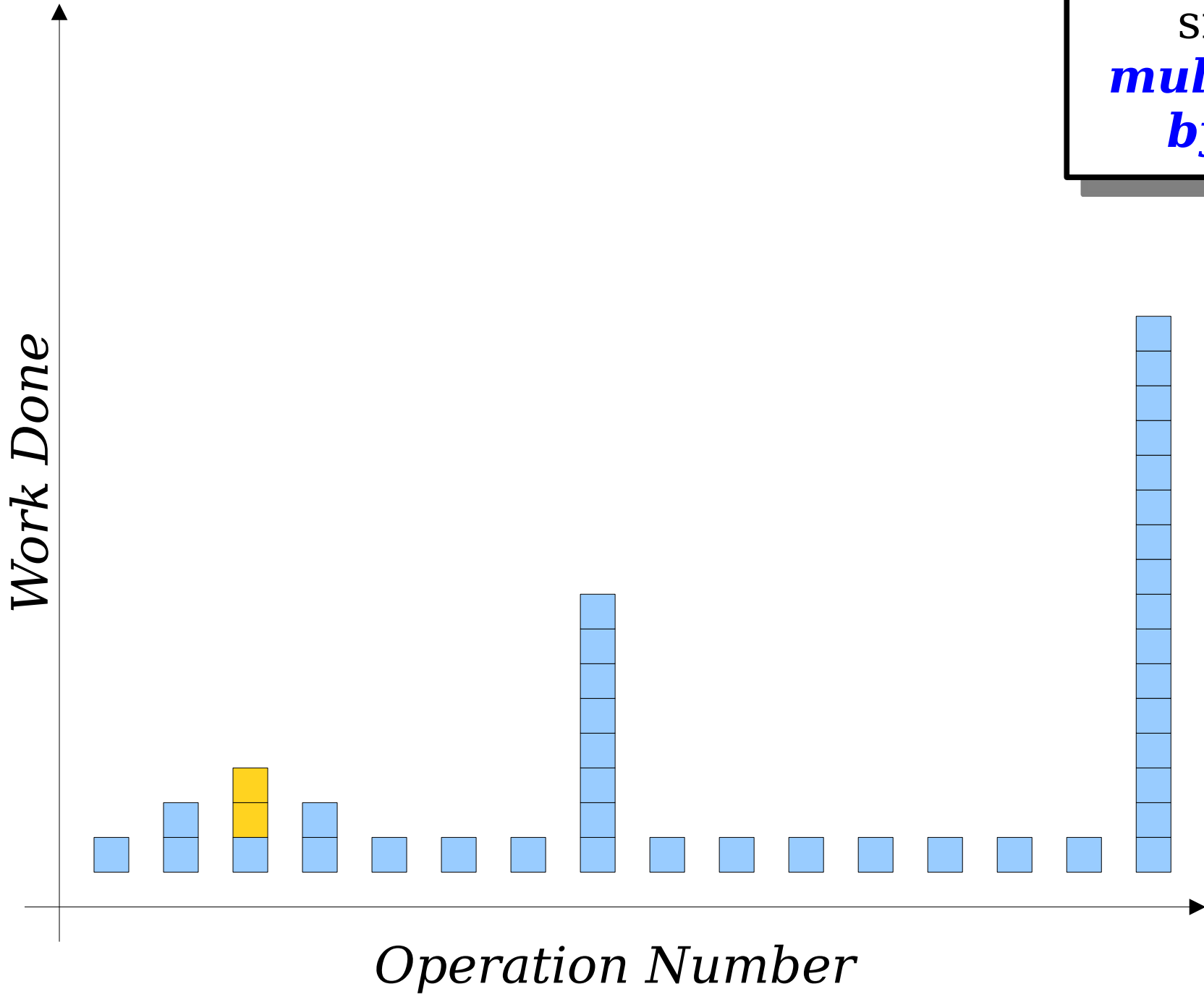
What's the average work done with each push?  
To find out, let's see how much total work was done.



Increase array size by *multiplying by two.*

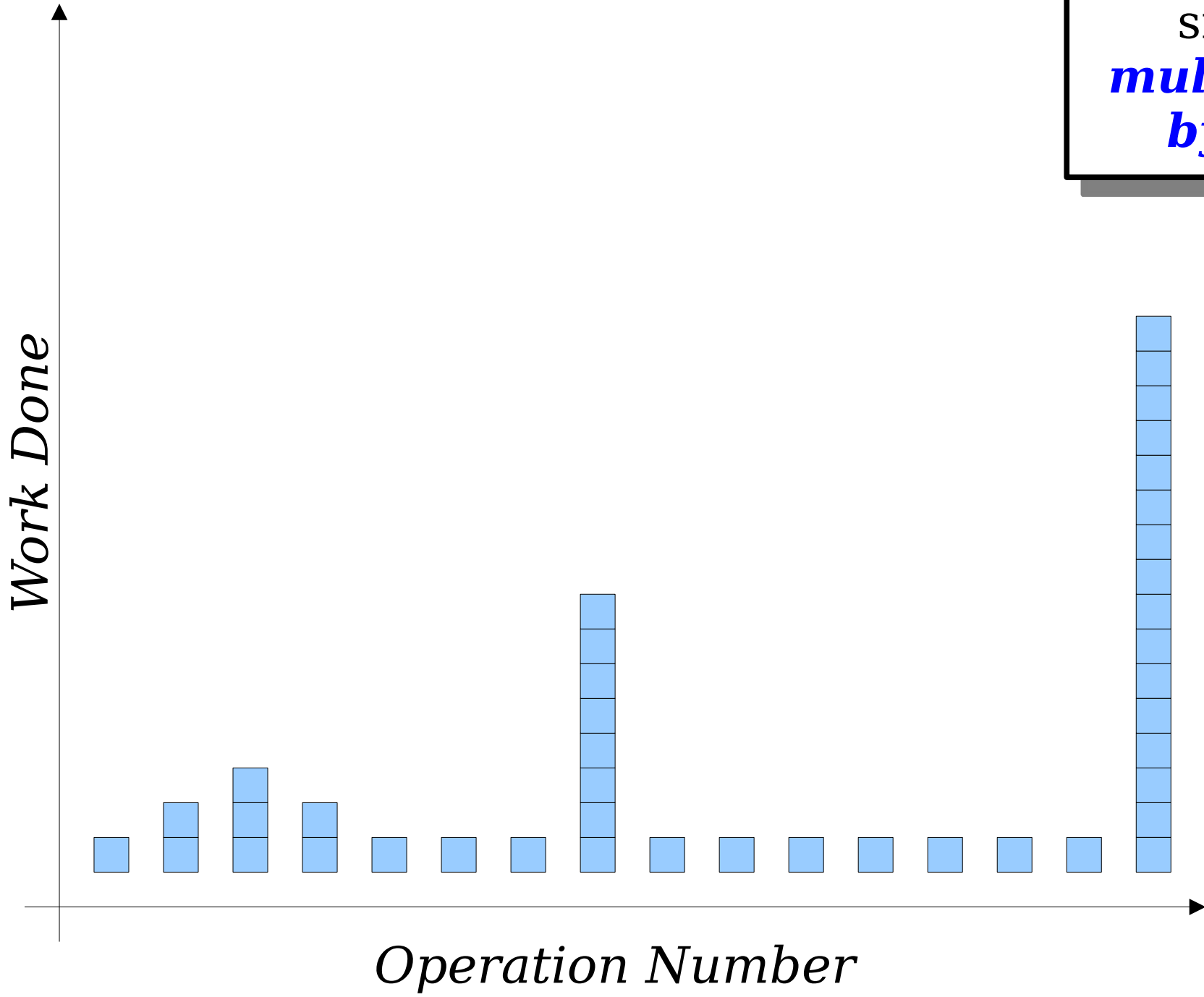


Increase array size by *multiplying by two.*

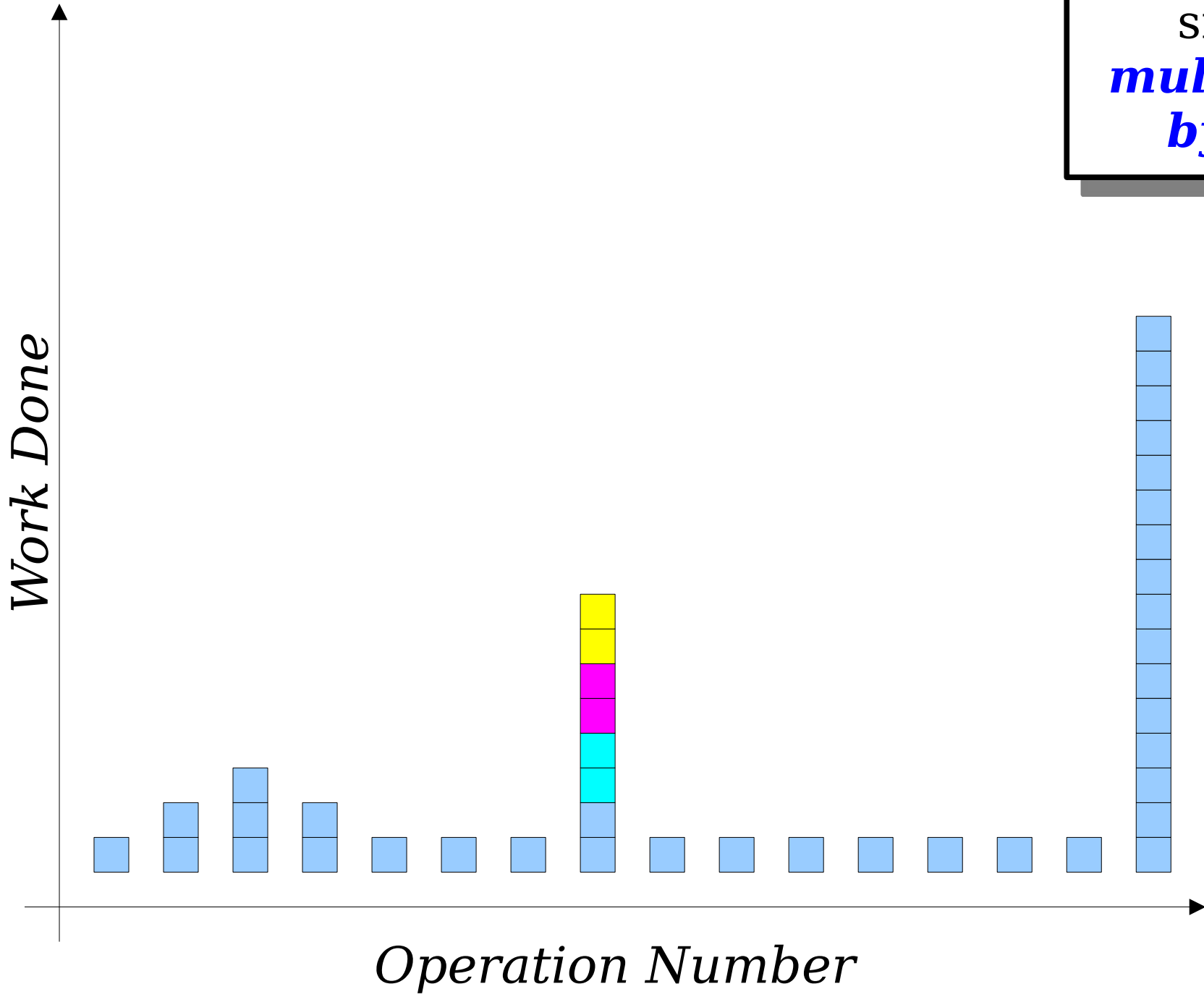




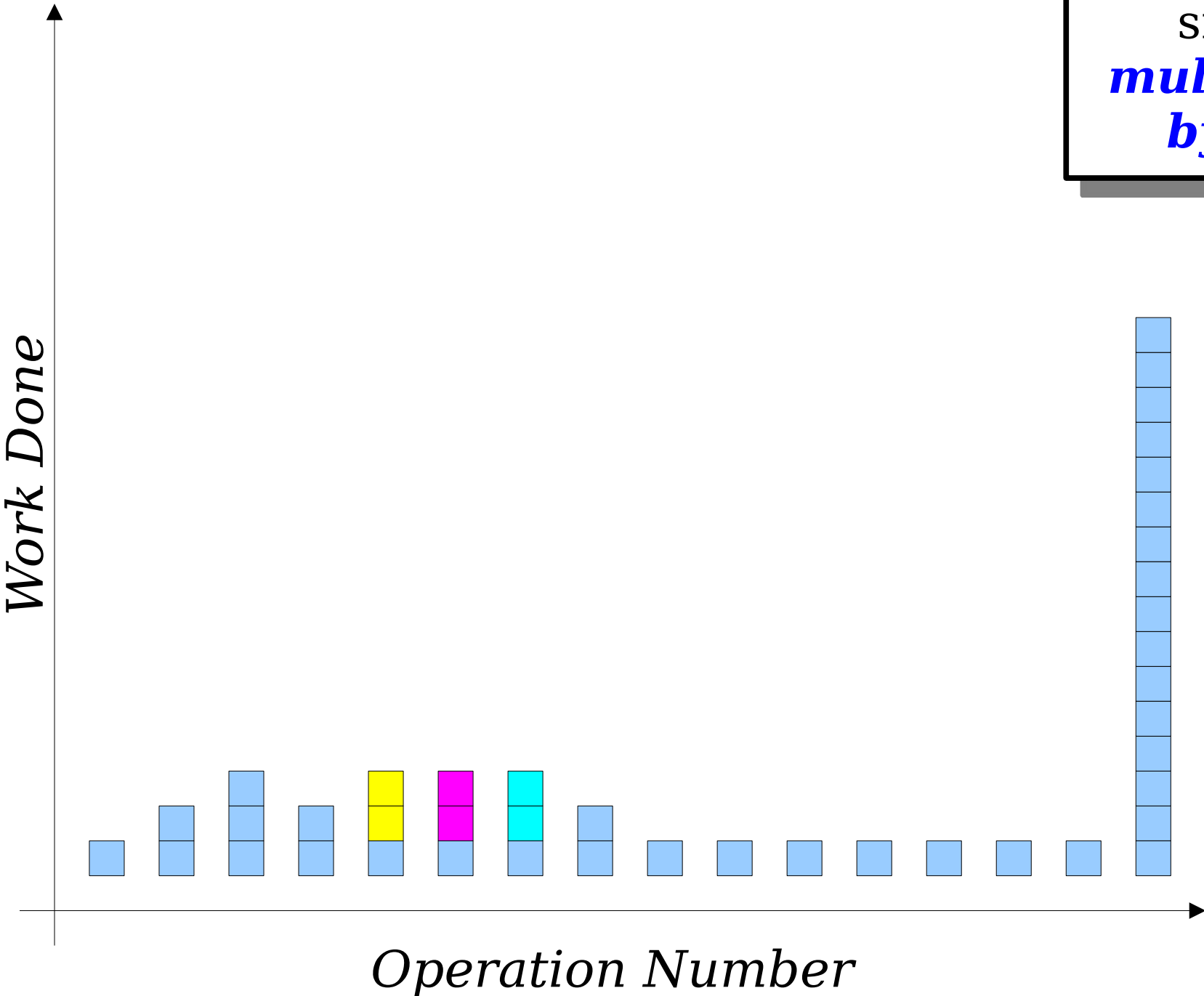
Increase array size by *multiplying by two.*



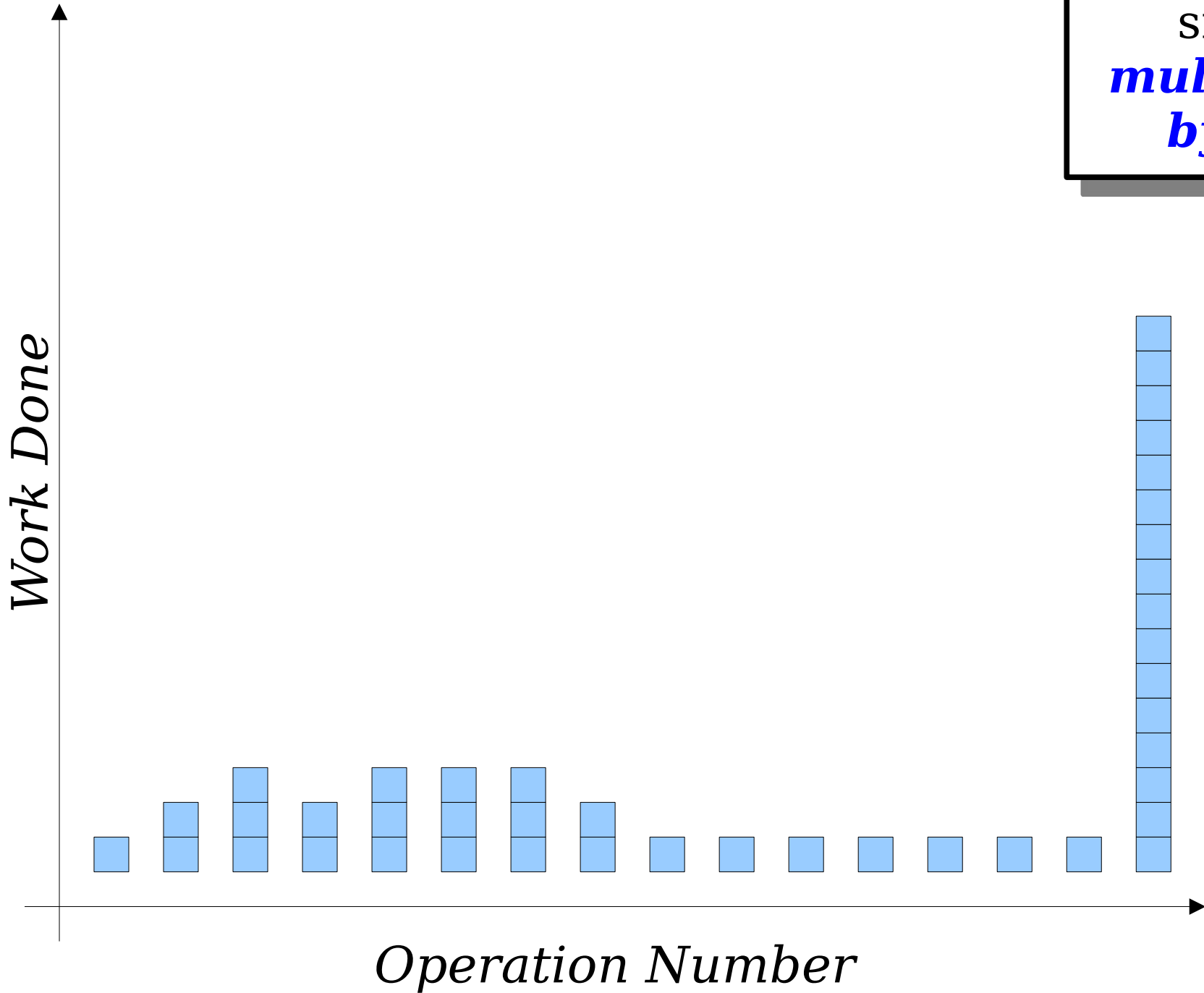
Increase array size by *multiplying by two.*



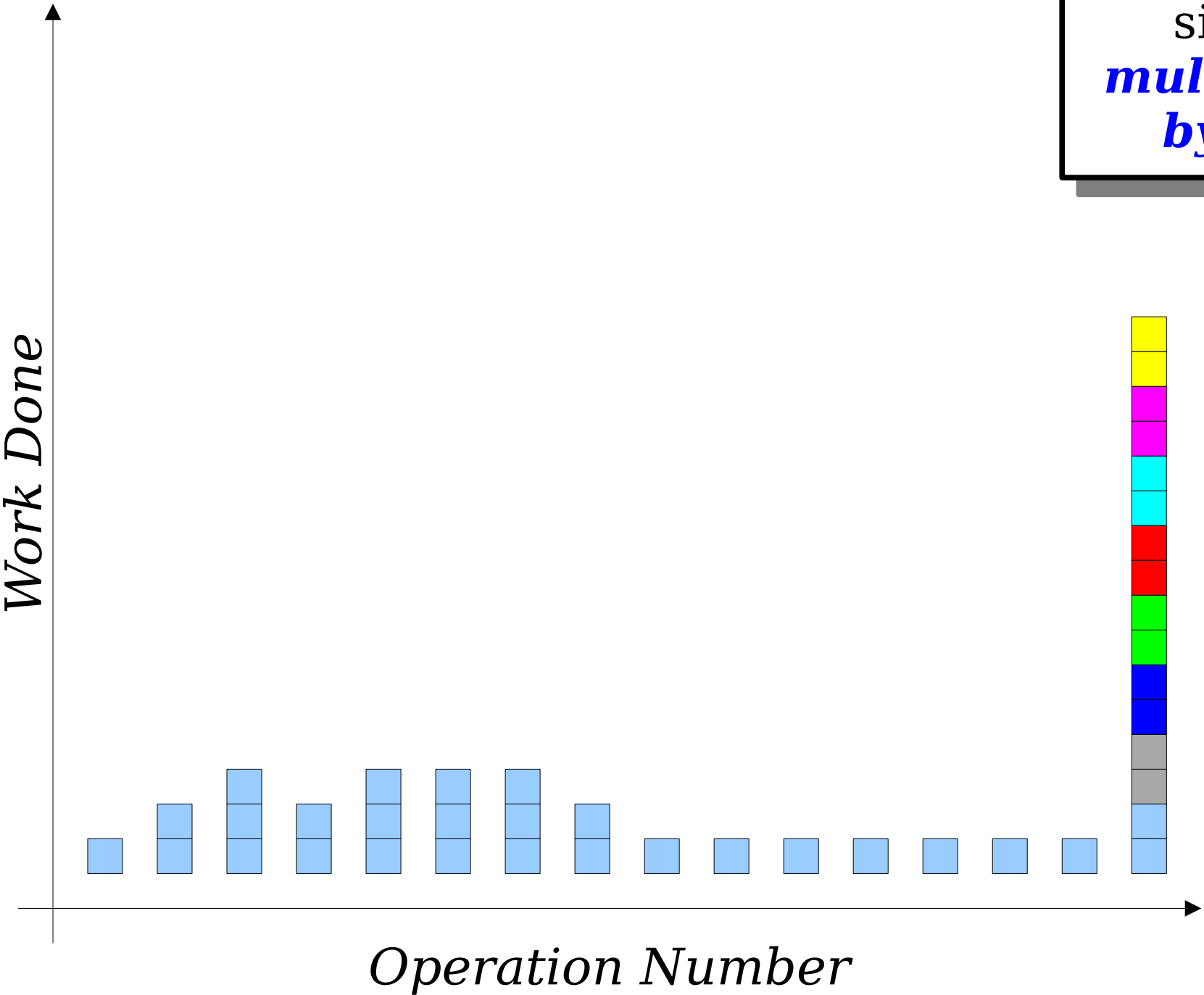
Increase array size by *multiplying by two.*



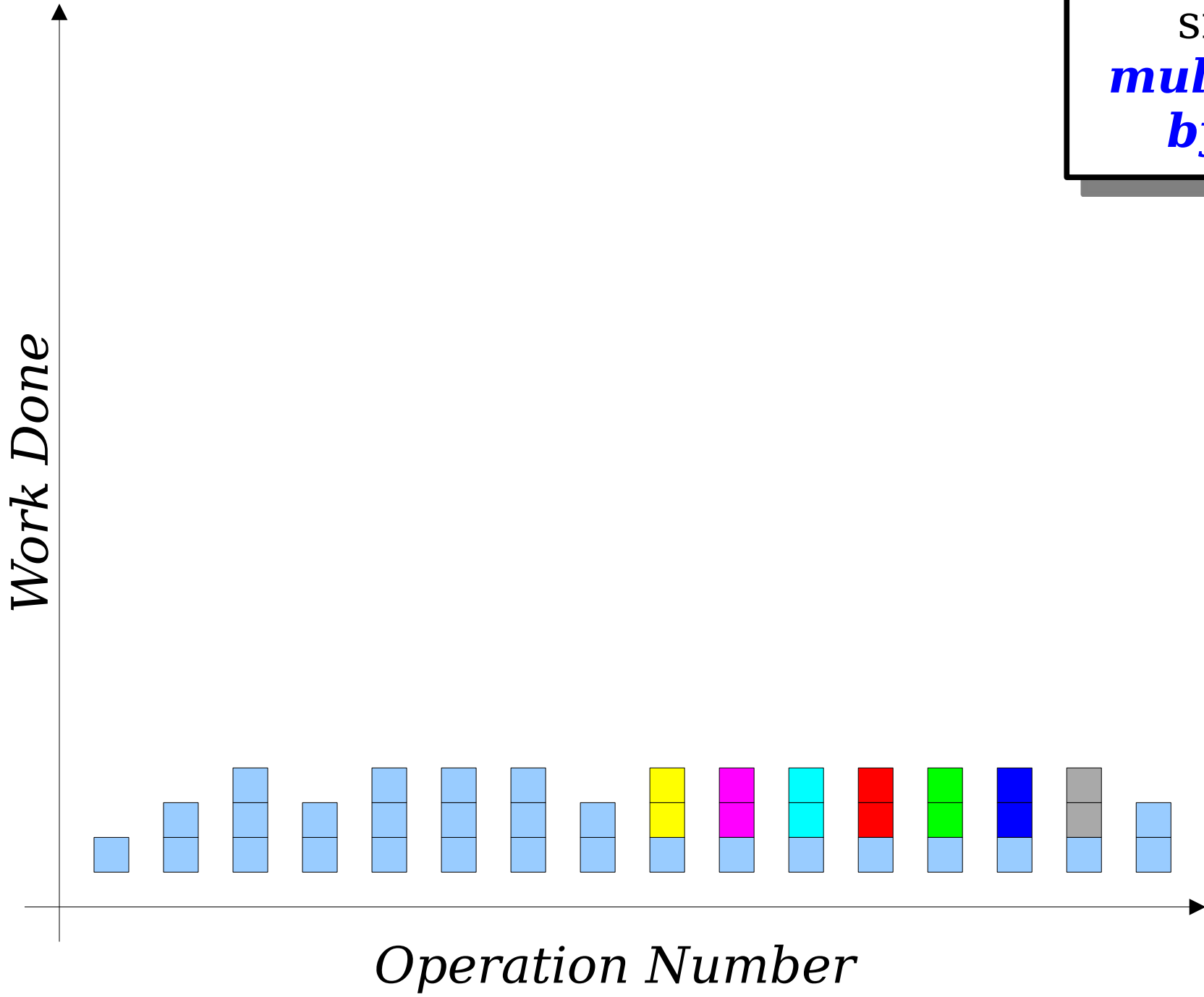
Increase array size by *multiplying by two.*



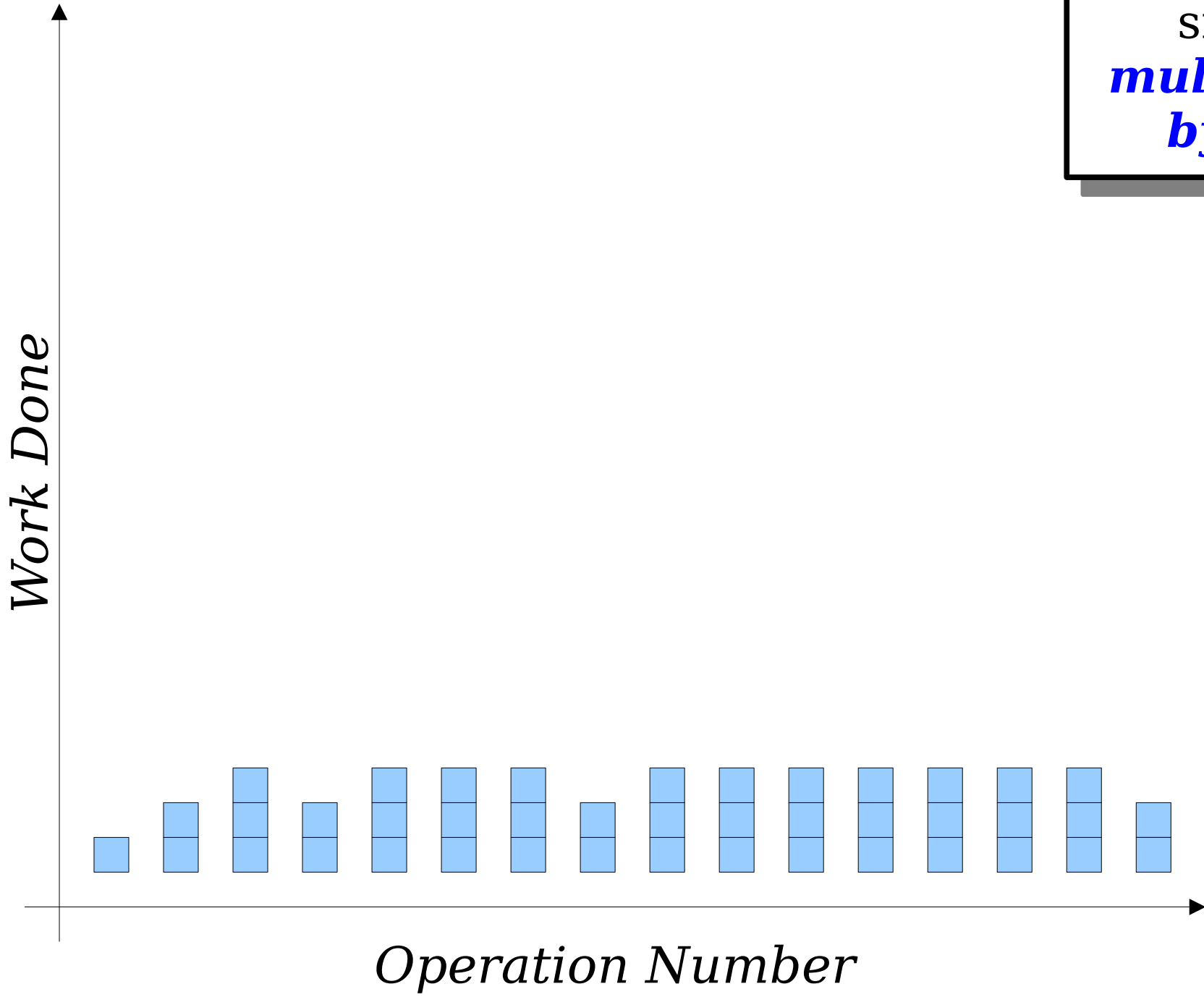
Increase array size by *multiplying by two*.



Increase array size by *multiplying by two.*

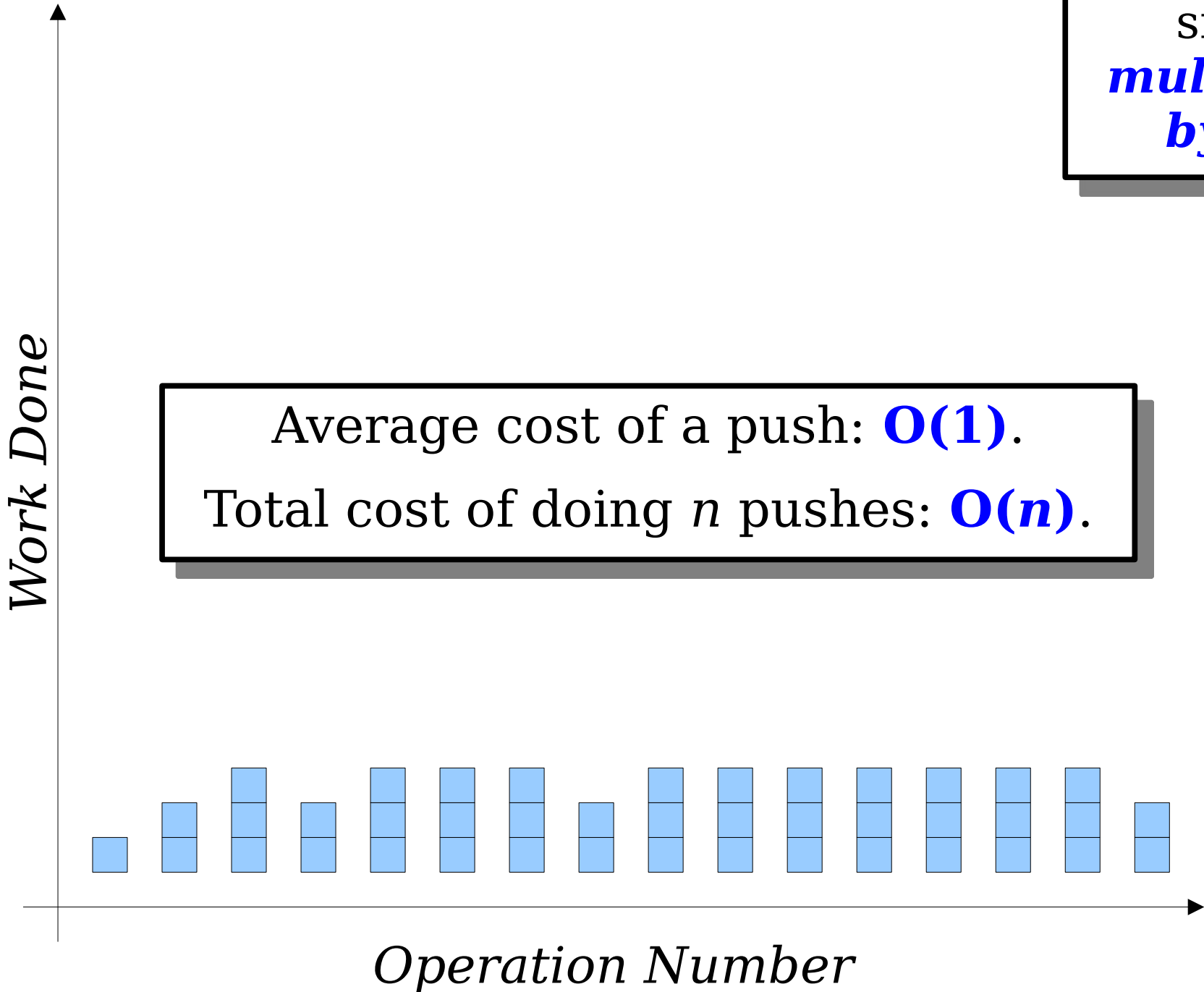


Increase array size by  
size by  
***multiplying  
by two.***



Increase array size by *multiplying by two*.

Average cost of a push:  $O(1)$ .  
Total cost of doing  $n$  pushes:  $O(n)$ .





# Amortized Analysis

- The analysis we have just done is called an *amortized analysis*.
- We reason about the total work done by allowing ourselves to backcharge work to previous operations, then look at the “average” amount of work done per operation.
- In an amortized sense, our implementation of the stack is extremely fast!
- This is one of the most common approaches to implementing Stack (and Vector, for that matter).

# Summary for Today

- We can make our stack grow by creating new arrays any time we run out of space.
- Growing that array by one extra slot or two extra slots uses little memory, but makes pushes expensive (average cost  $O(n)$ ).
- Doubling the size of the array when we run out of space uses more memory, but makes pushes cheap (amortized cost  $O(1)$ ).
- In practice, it's worth paying this slight space cost for a marked improvement in runtime.

# Your Action Items

- ***Start Assignment 6.***
  - Slow and steady progress is the name of the game here.
  - Ask for help if you need it! That's what we're here for.
- ***Review your midterm feedback.***
  - Regrade requests are due on Monday, so make sure you've read over the graders' comments by then.

# Next Time

- ***Hash Functions***
  - A magical and wonderful gift from the world of mathematics.
- ***Hash Tables***
  - How do we implement Map and Set?