

Programming Abstractions (Accelerated)
Winter 2017
Stanford University
Computer Science Department

THE LIFE CHANGING MAGIC OF

DIJKSTRA AND A*

Friday, March 10, 2017

Reading: Programming Abstractions in C++, Chapter 18.6

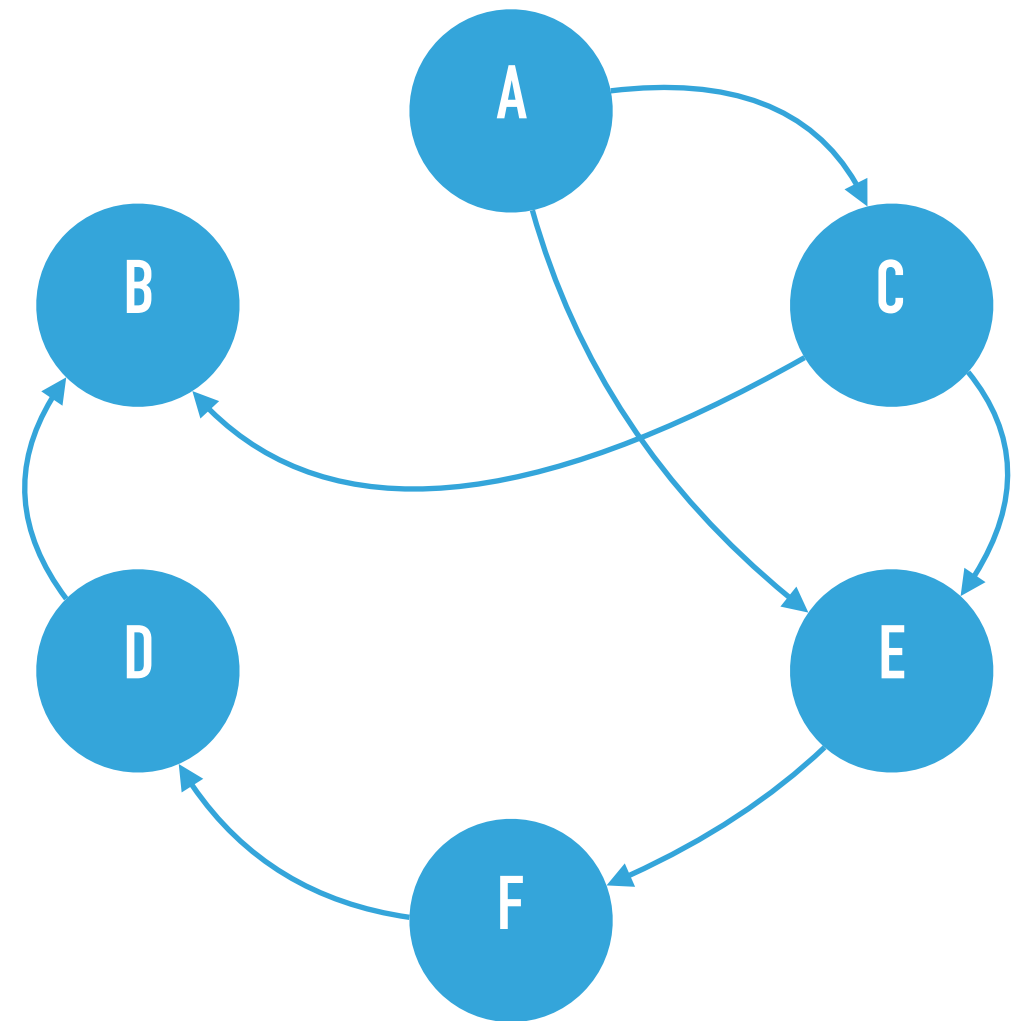
TODAY'S TOPICS – MORE GRAPHS!

- ▶ Reviewing DFS and BFS
- ▶ Comparing DFS and BFS
- ▶ Making weighty decisions using Dijkstra's algorithm
- ▶ Looking into the future with A*
- ▶ Google Maps

REVIEWING DFS AND BFS

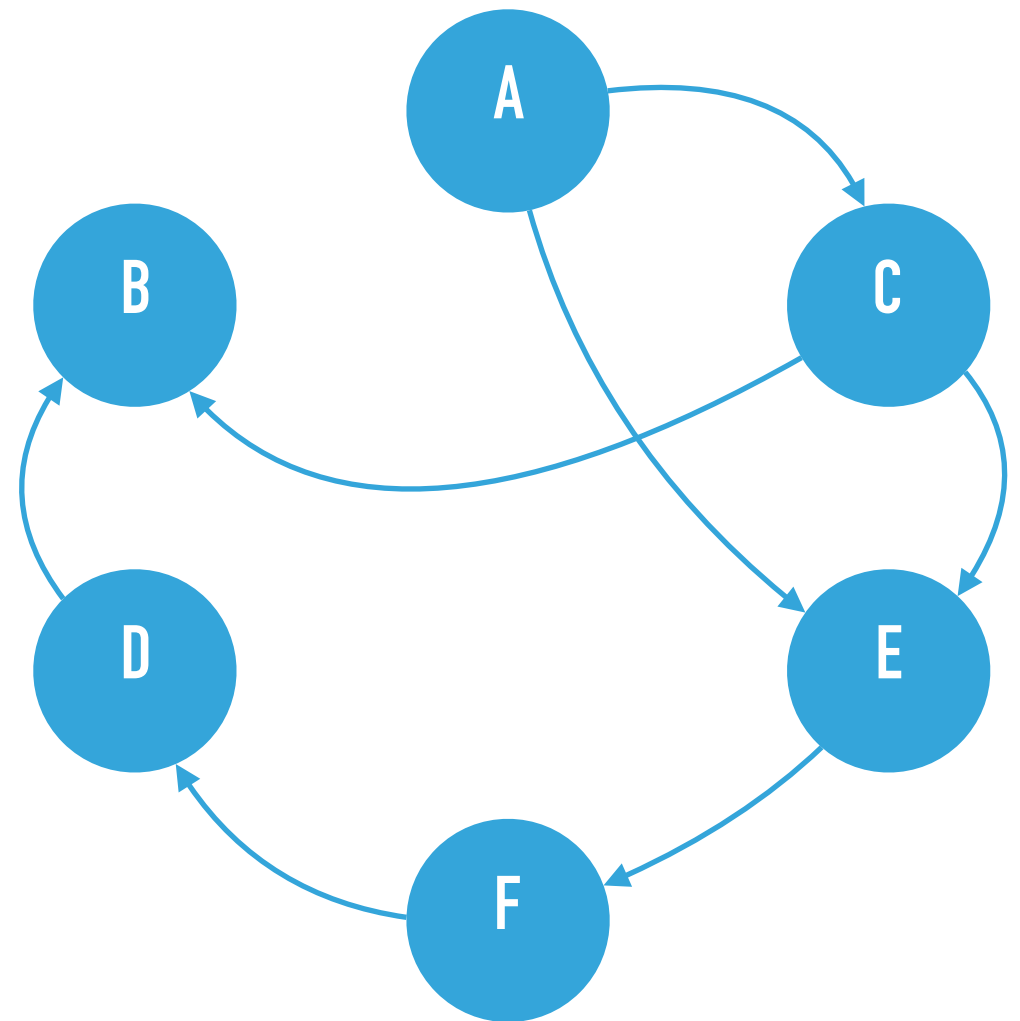
DEPTH FIRST SEARCH

- ▶ Find a path from A to B using *iterative* depth first search
- ▶ (Assume that nodes are pushed onto the stack in *alphabetic order*)



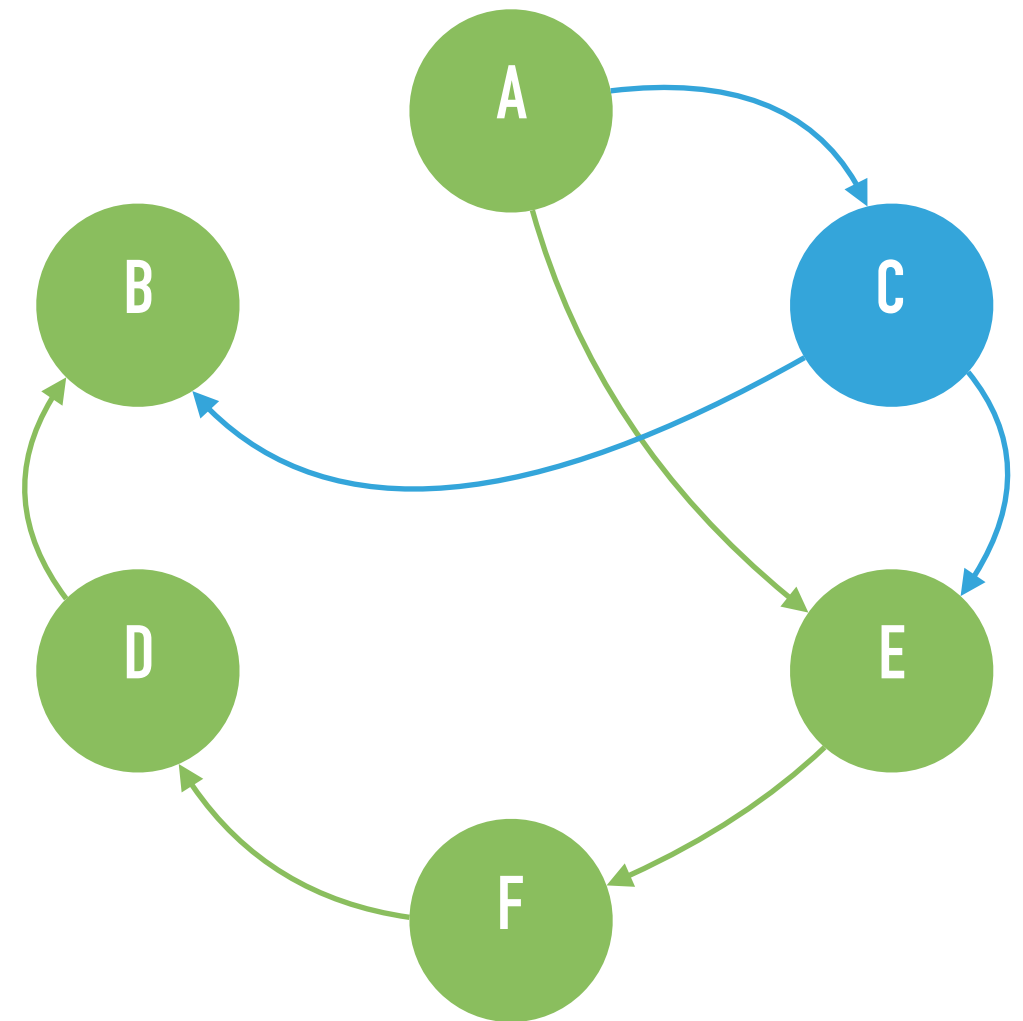
DEPTH FIRST SEARCH (ITERATIVE PSEUDOCODE)

- ▶ create a path with just start node and push onto stack s
- ▶ while s is not empty
 - ▶ $p = s.pop()$
 - ▶ $v = \text{last node of } p$
 - ▶ if v is end, you're done
 - ▶ mark v as visited
 - ▶ for each unvisited neighbor:
 - ▶ create new path and append neighbor
 - ▶ push new path onto s



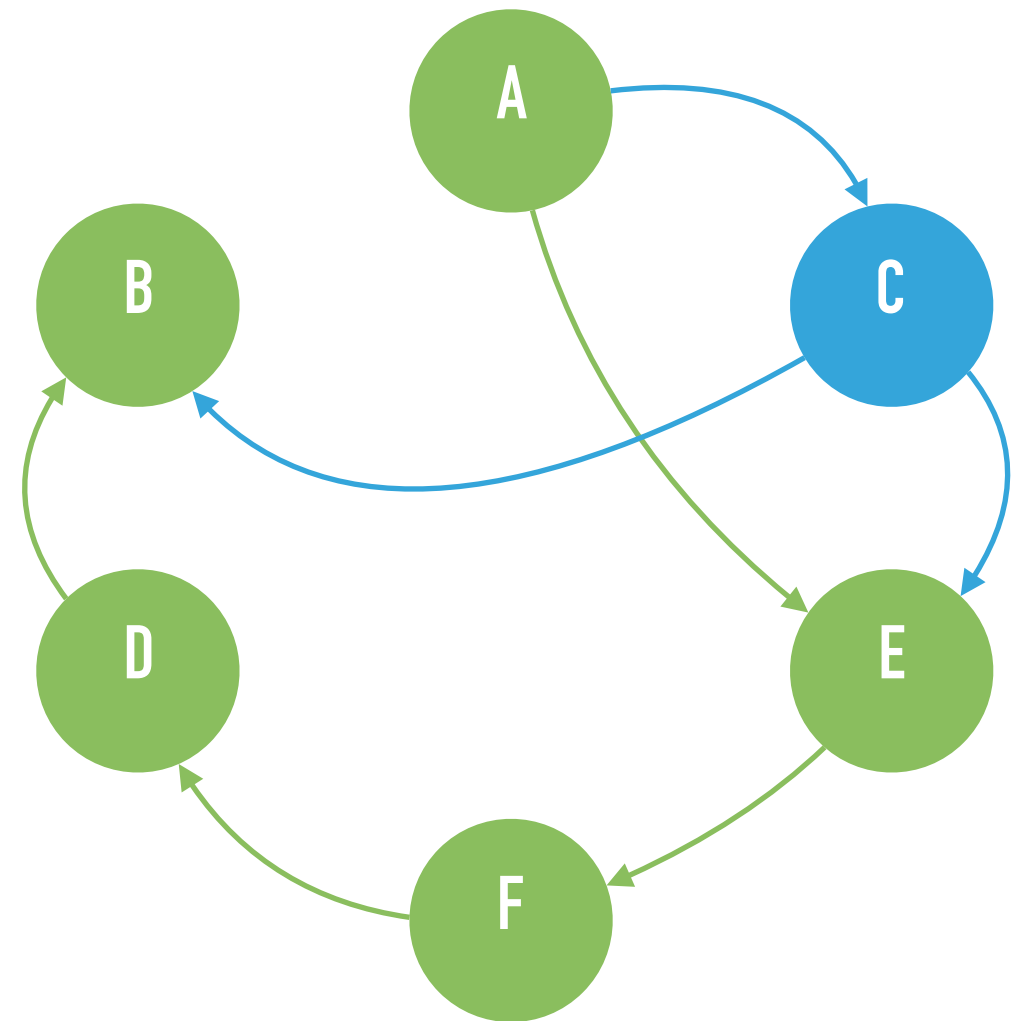
DEPTH FIRST SEARCH

- ▶ Find a path from A to B using *iterative* depth first search
 - ▶ (Assume that nodes are pushed onto the stack in *alphabetic order*)
- ▶ $A \rightarrow E \rightarrow F \rightarrow D \rightarrow B$

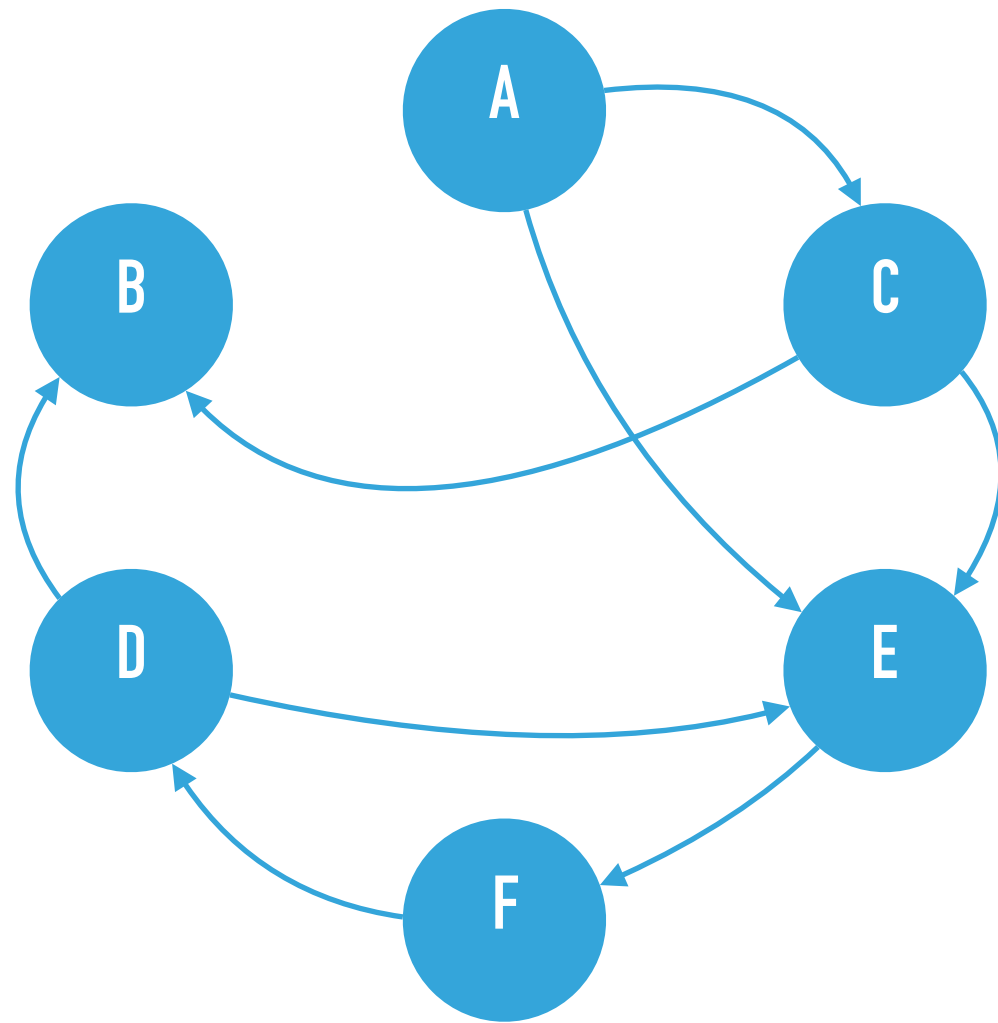


DEPTH FIRST SEARCH

- ▶ Find a path from A to B using *iterative* depth first search
 - ▶ (Assume that nodes are pushed onto the stack in *alphabetic order*)
- ▶ $A \rightarrow E \rightarrow F \rightarrow D \rightarrow B$
- ▶ Is this the shortest path?



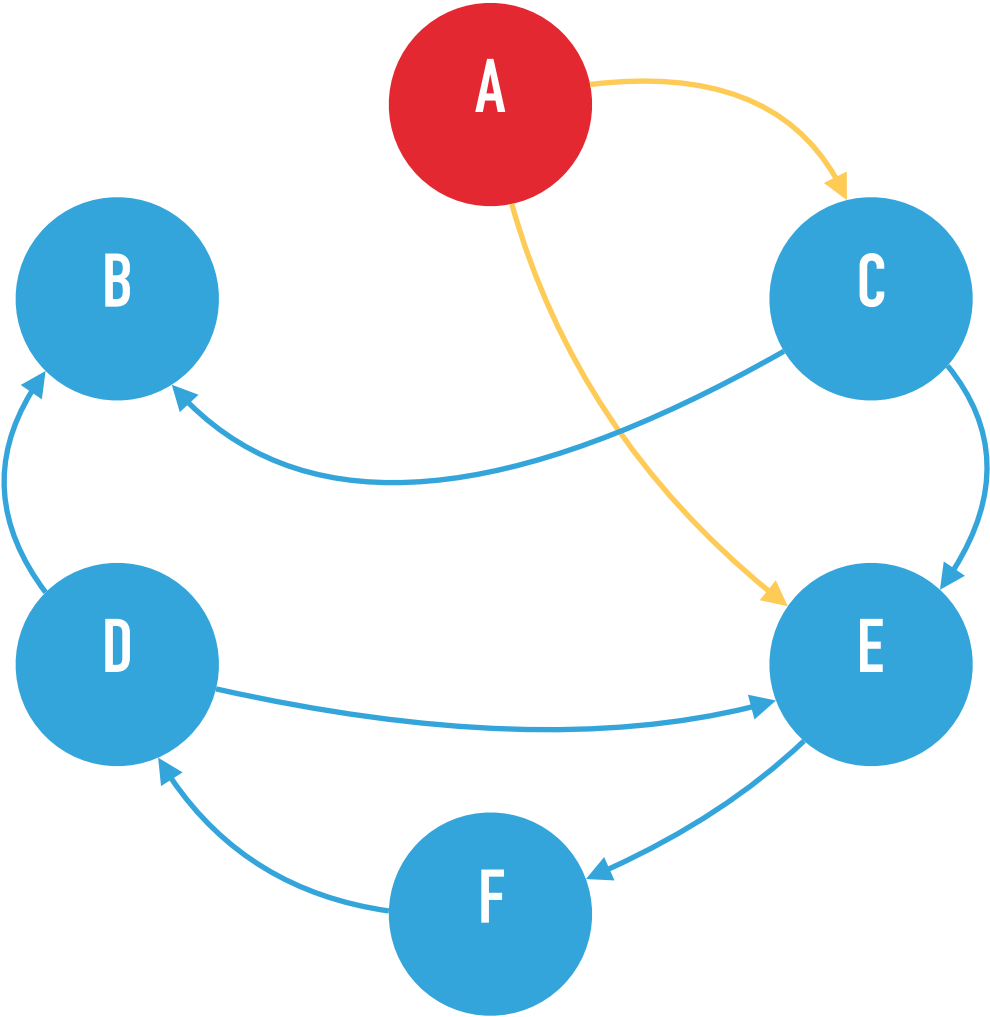
DEPTH FIRST SEARCH



Paths to Consider (Stack)



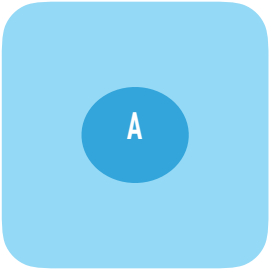
DEPTH FIRST SEARCH



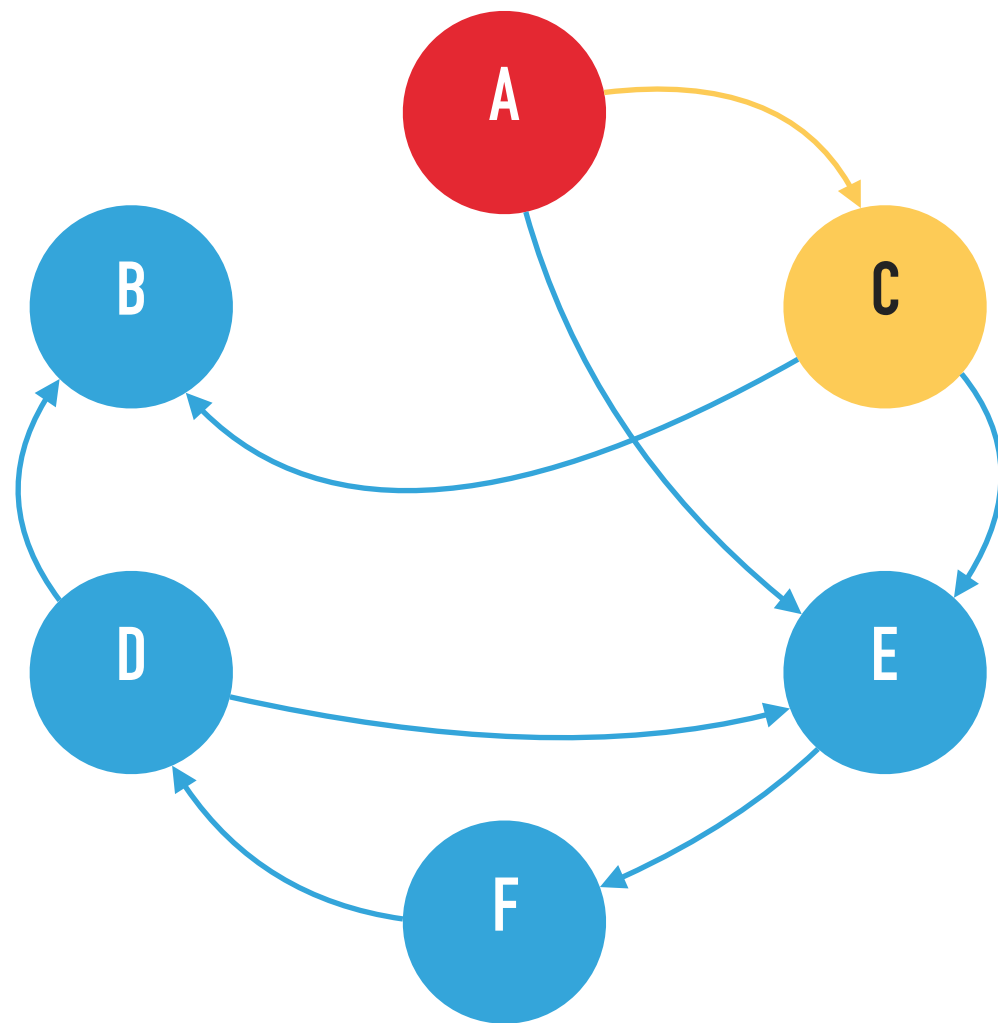
Paths to Consider (Stack)



Current Path



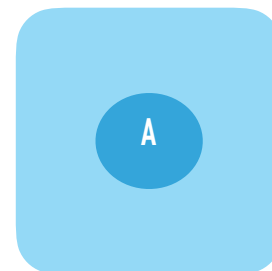
DEPTH FIRST SEARCH



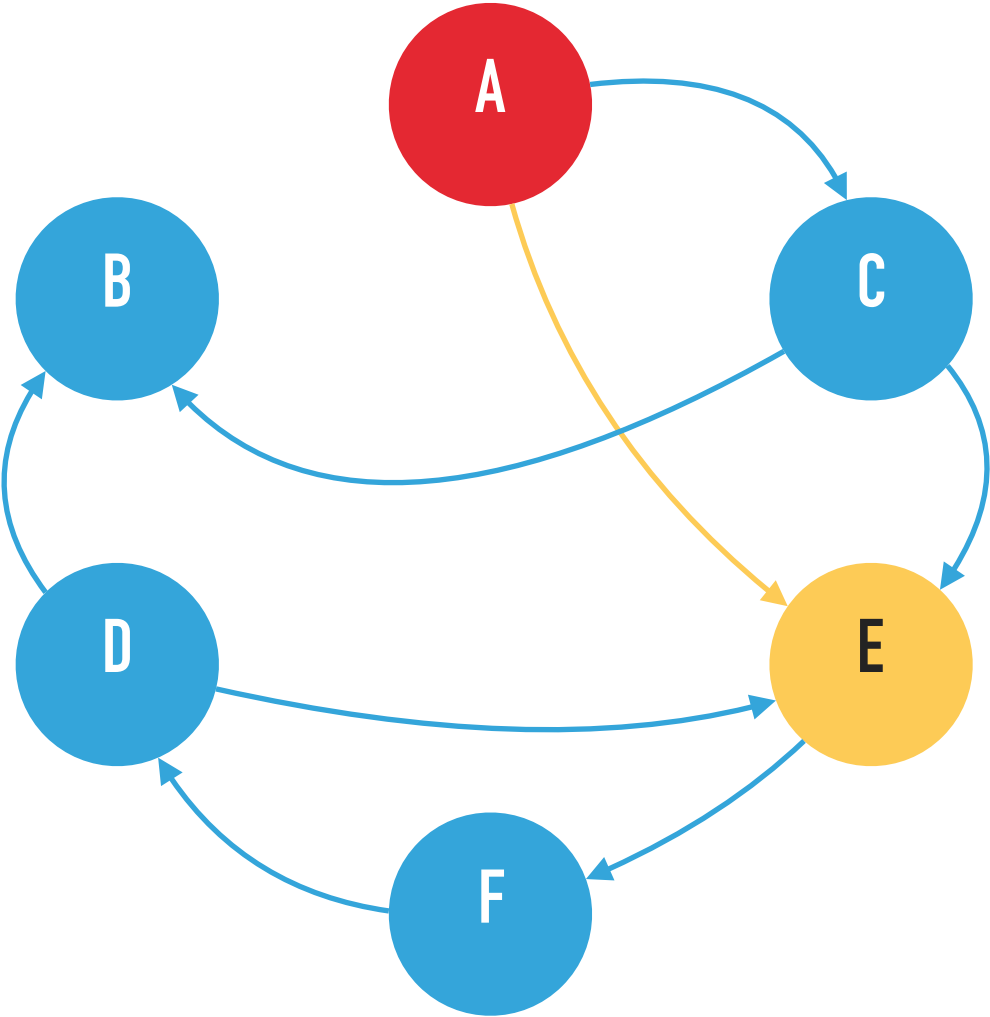
Paths to Consider (Stack)



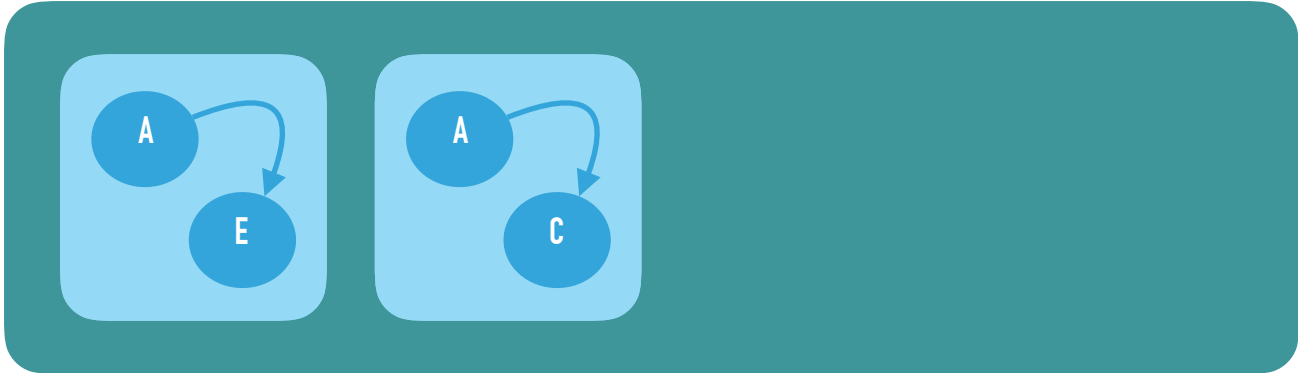
Current Path



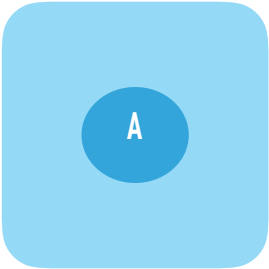
DEPTH FIRST SEARCH



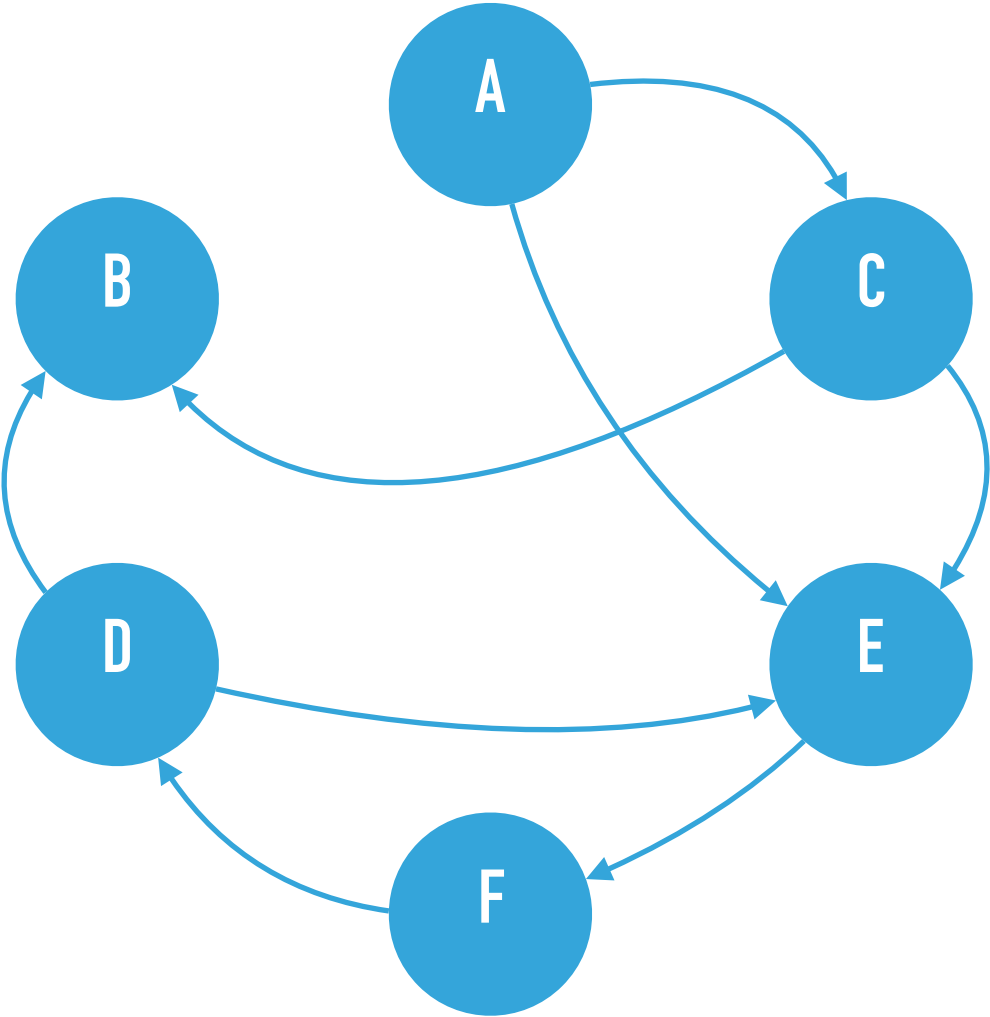
Paths to Consider (Stack)



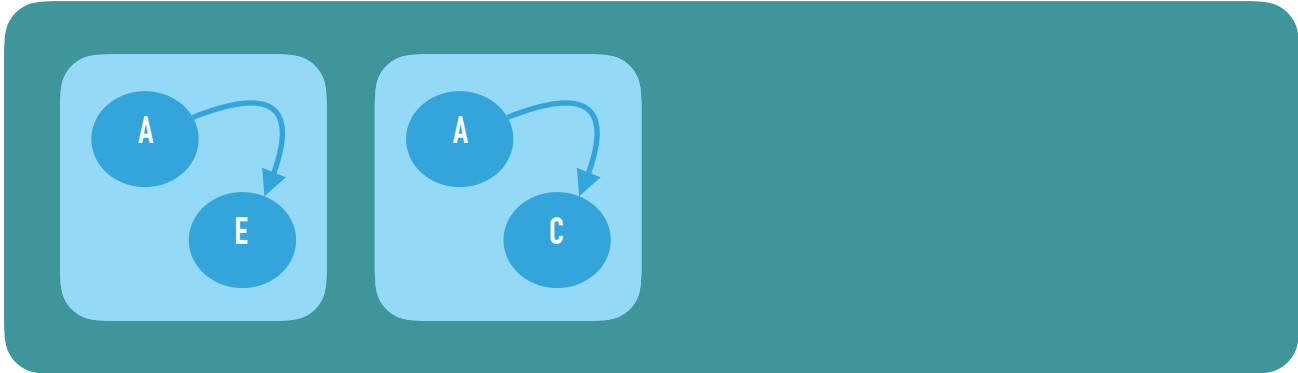
Current Path



DEPTH FIRST SEARCH

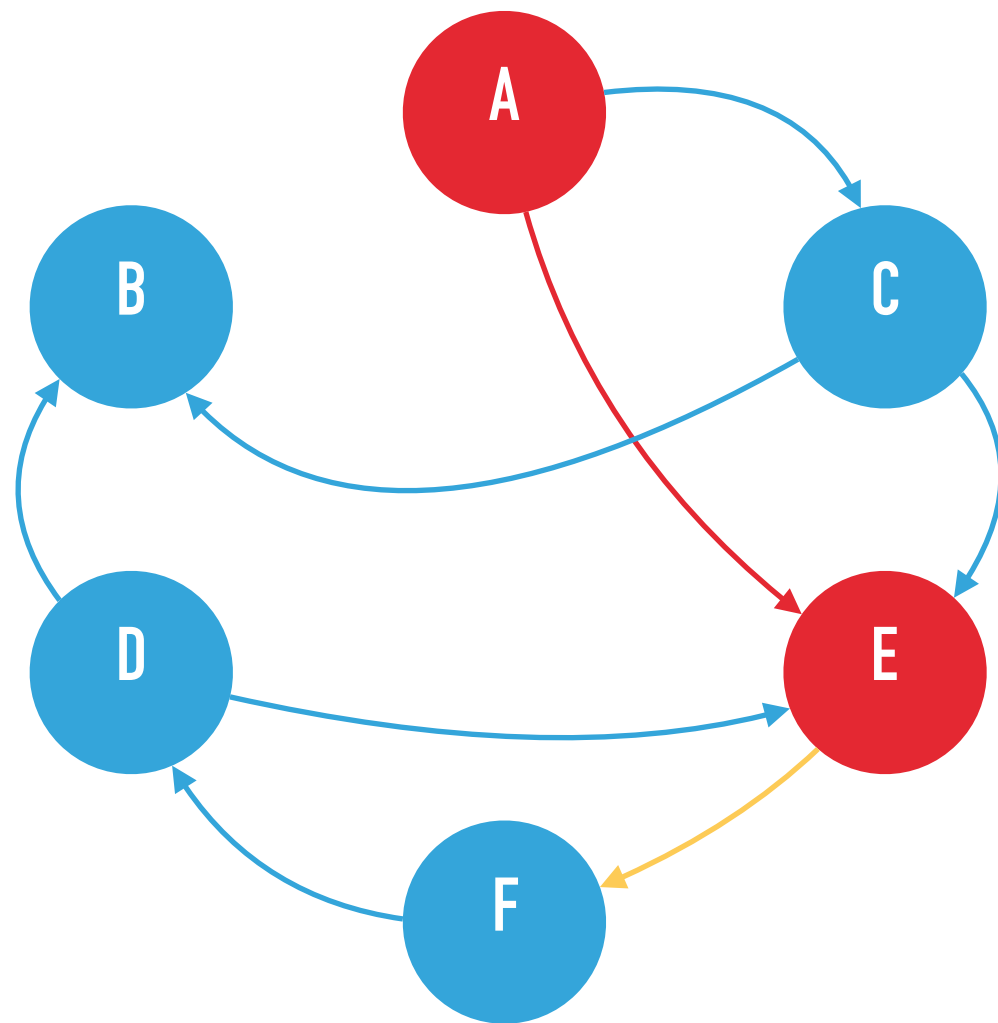


Paths to Consider (Stack)



Current Path

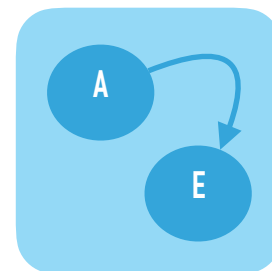
DEPTH FIRST SEARCH



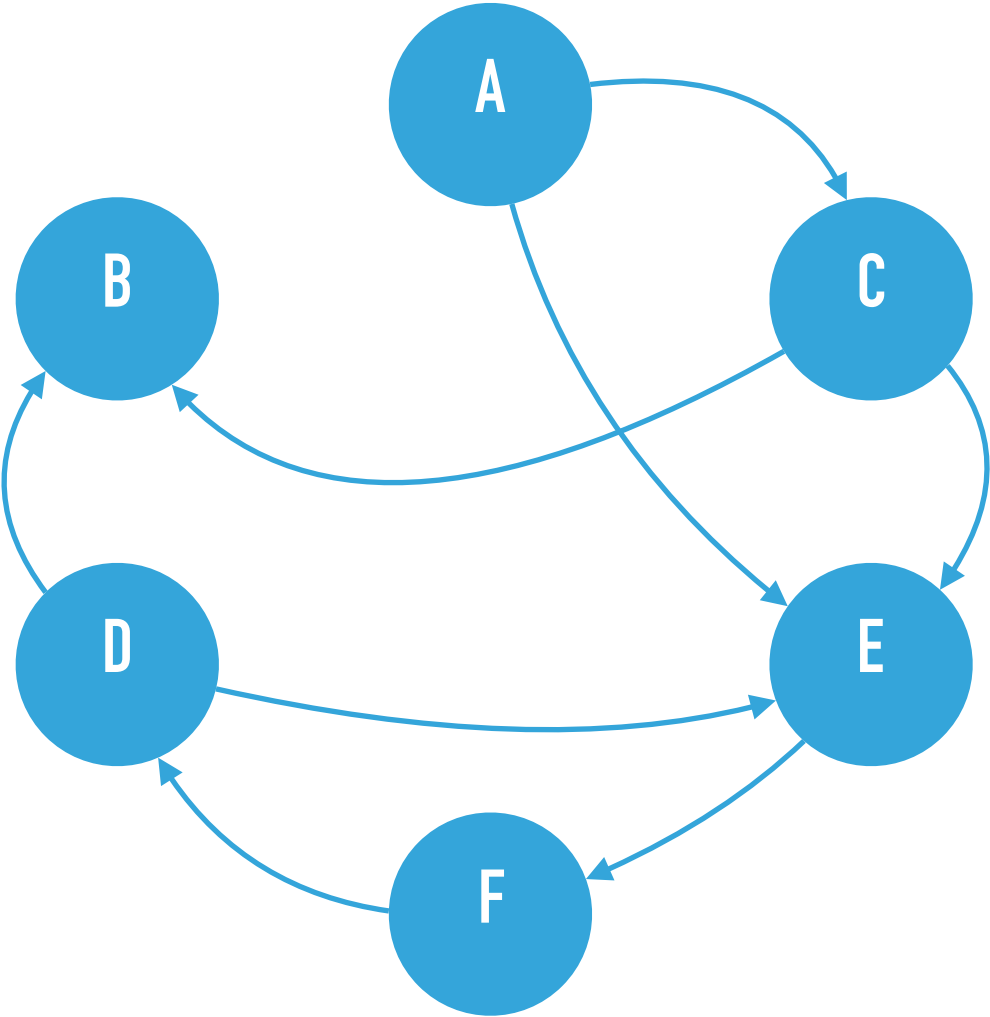
Paths to Consider (Stack)



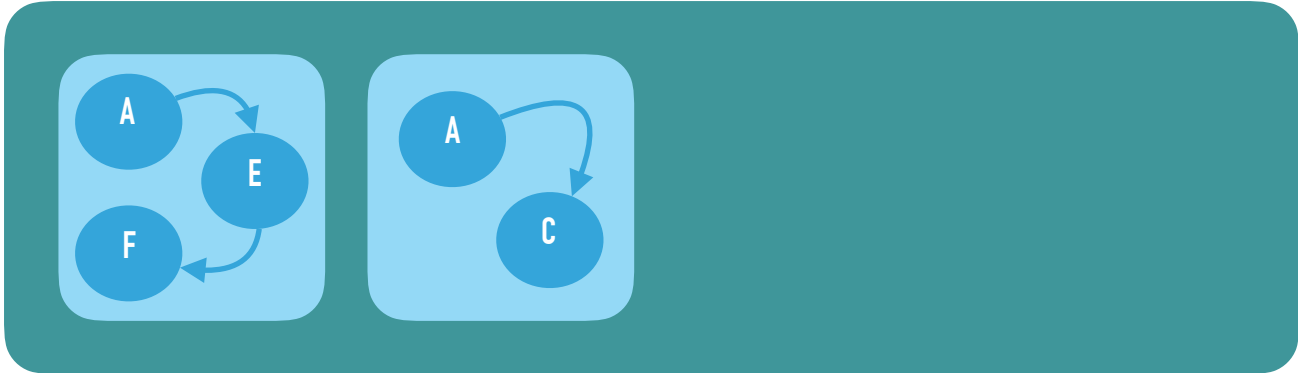
Current Path



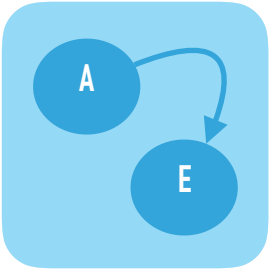
DEPTH FIRST SEARCH



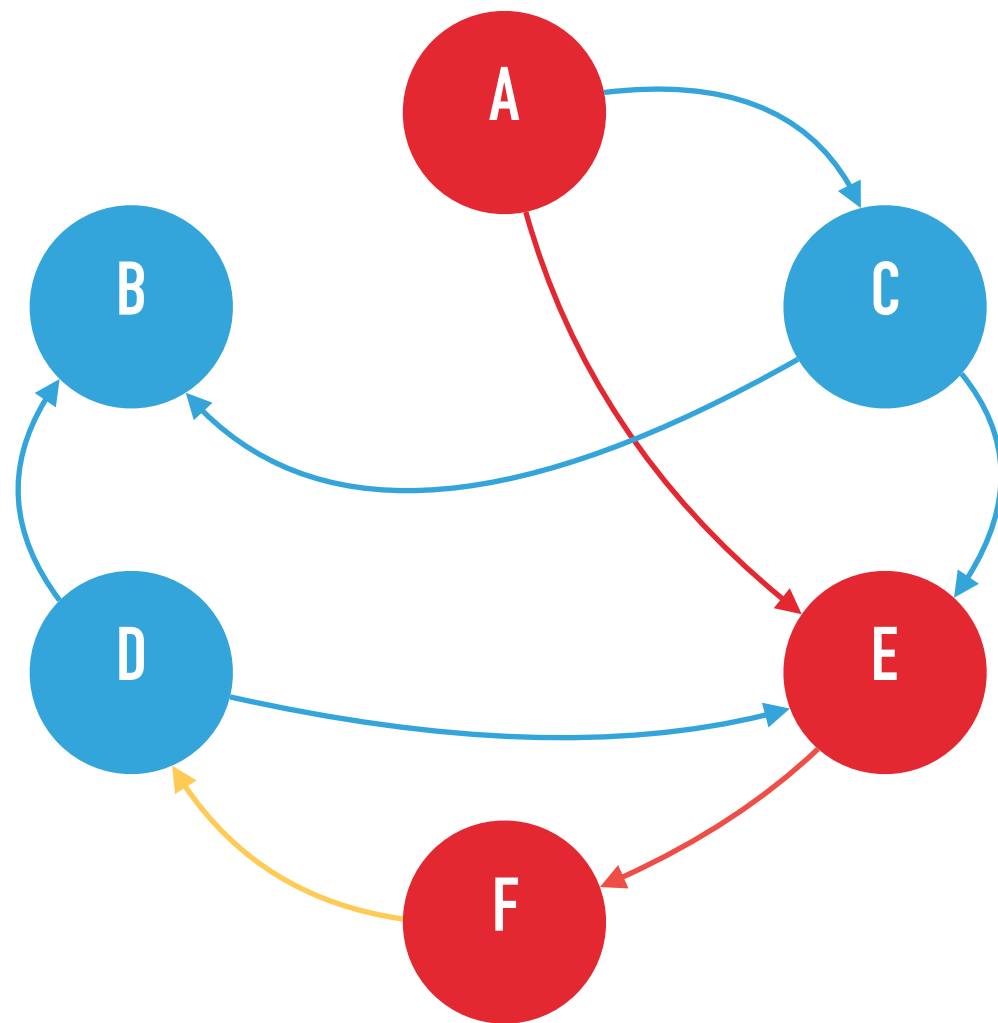
Paths to Consider (Stack)



Current Path



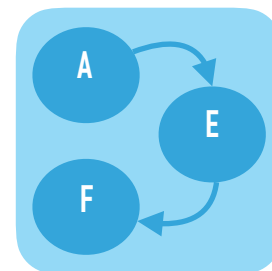
DEPTH FIRST SEARCH



Paths to Consider (Stack)

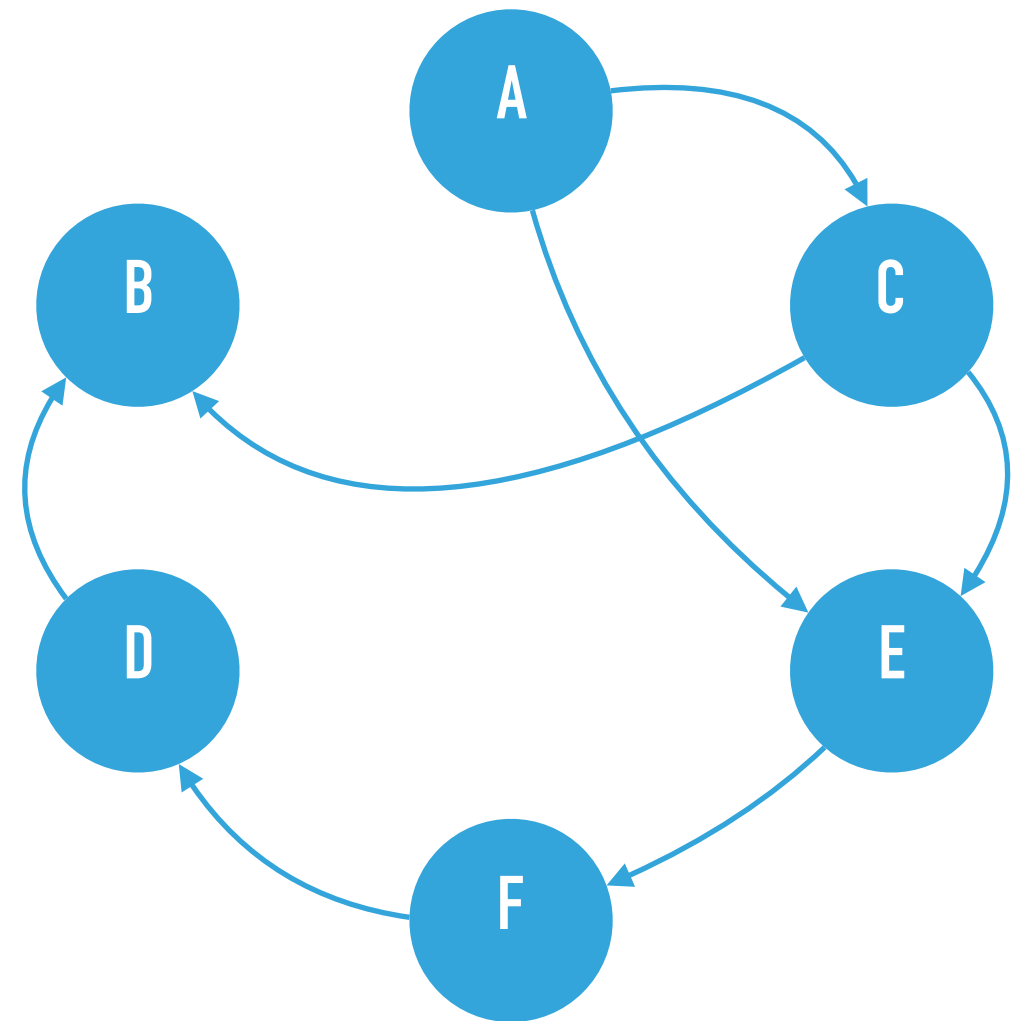


Current Path



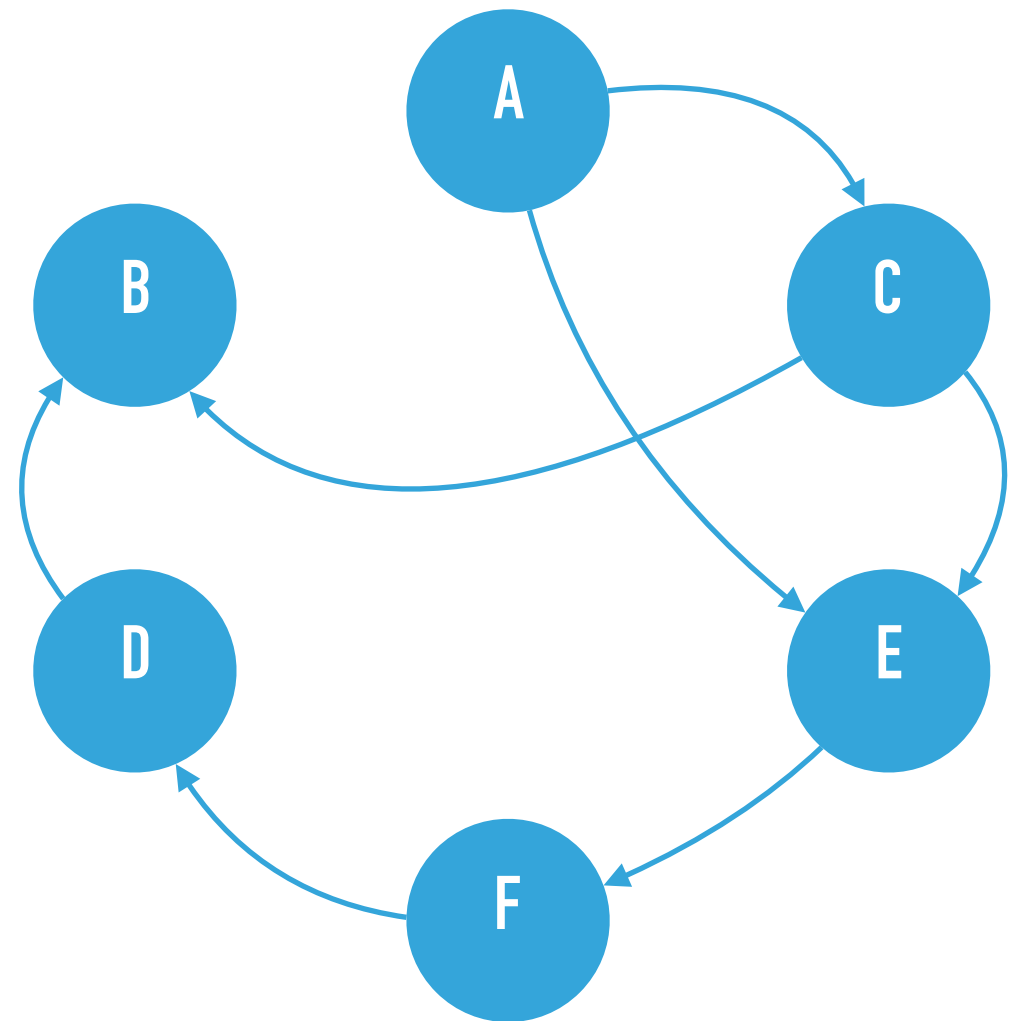
BREADTH FIRST SEARCH

- ▶ Find a path from A to B using breadth first search
- ▶ (Assume that nodes are pushed onto the queue in *alphabetic order*)



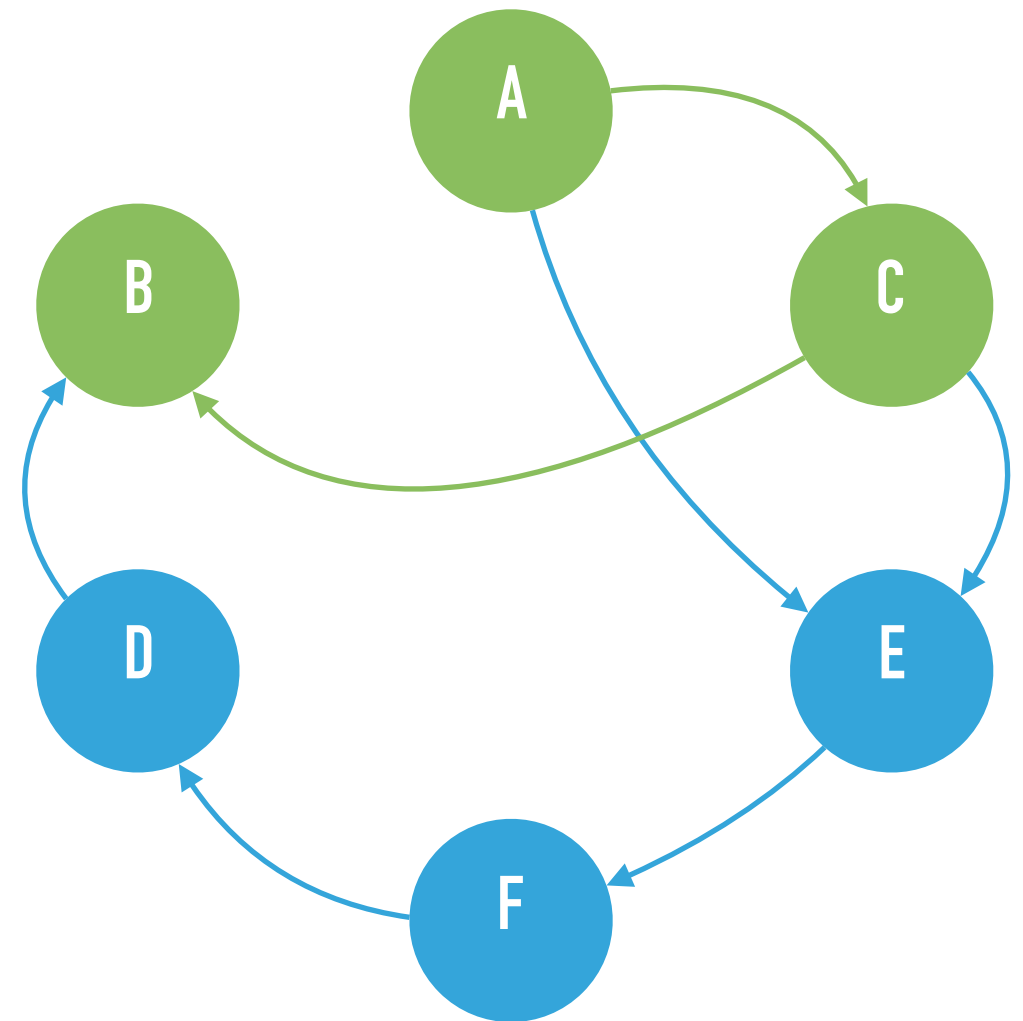
BREADTH FIRST SEARCH (PSEUDOCODE)

- ▶ create a path with just start node and enqueue into queue q
- ▶ while q is not empty
 - ▶ $p = q.dequeue()$
 - ▶ $v = \text{last node of } p$
 - ▶ if v is end, you're done
 - ▶ mark v as visited
 - ▶ for each unvisited neighbor:
 - ▶ create new path and append neighbor
 - ▶ enqueue new path into q



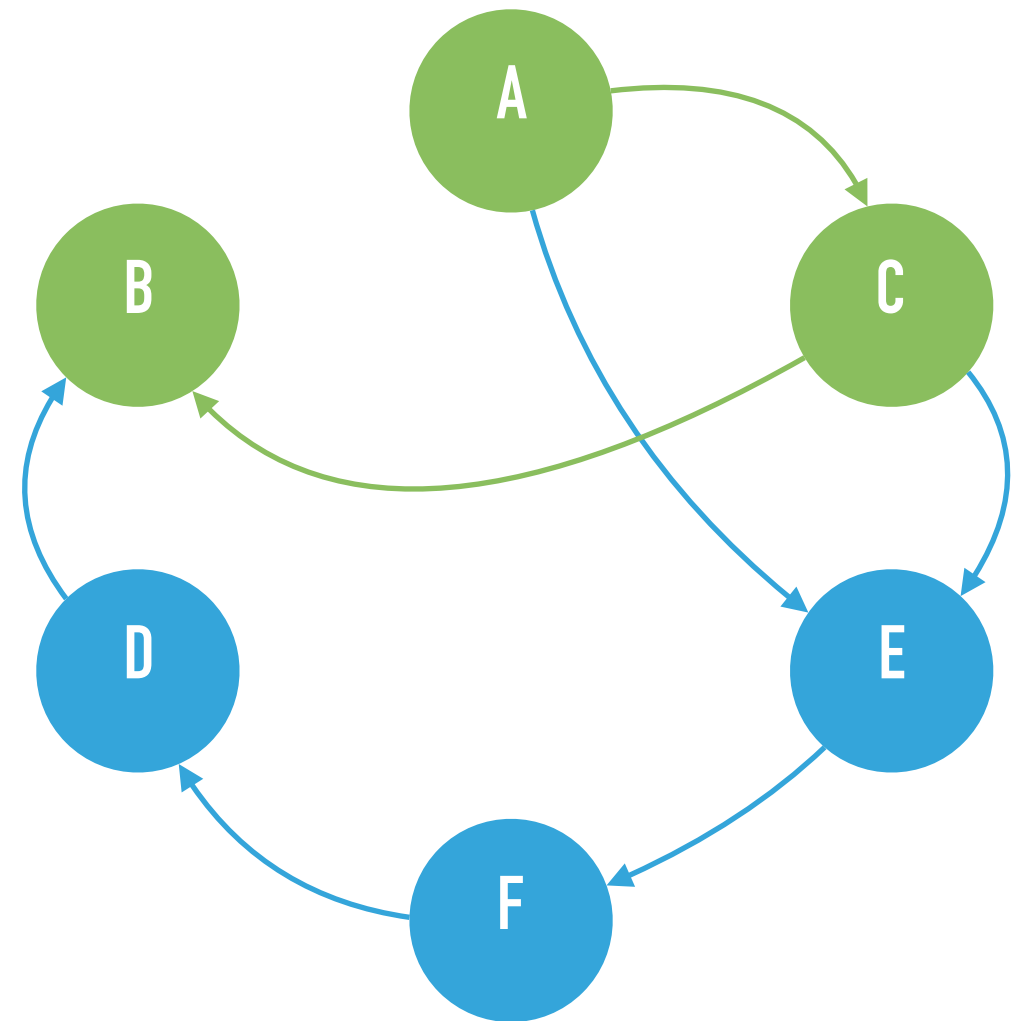
BREADTH FIRST SEARCH

- ▶ Find a path from A to F using breadth first search
 - ▶ (Assume that nodes are pushed onto the queue in *alphabetic order*)
- ▶ $A \rightarrow C \rightarrow B$



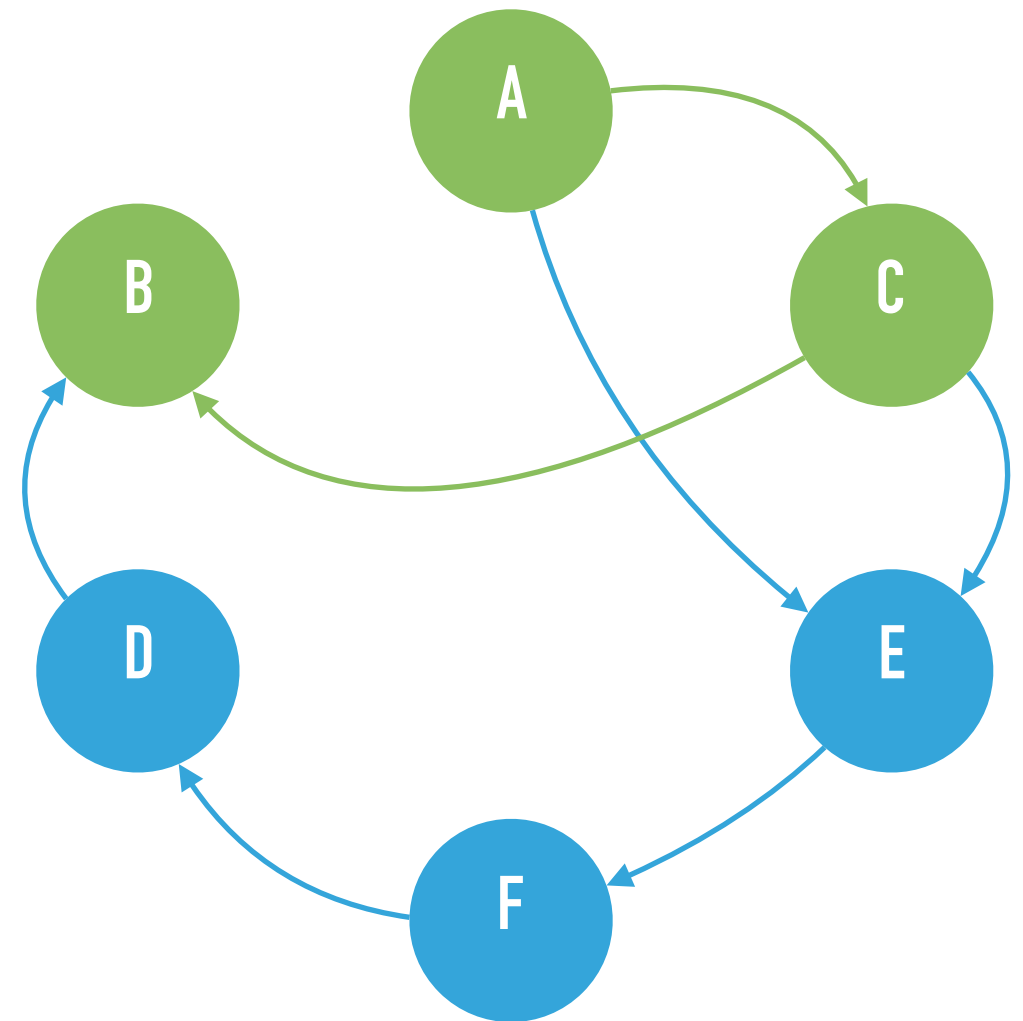
BREADTH FIRST SEARCH

- ▶ Find a path from A to F using breadth first search
 - ▶ (Assume that nodes are pushed onto the queue in *alphabetic order*)
- ▶ $A \rightarrow C \rightarrow B$
- ▶ Is *this* the shortest path?

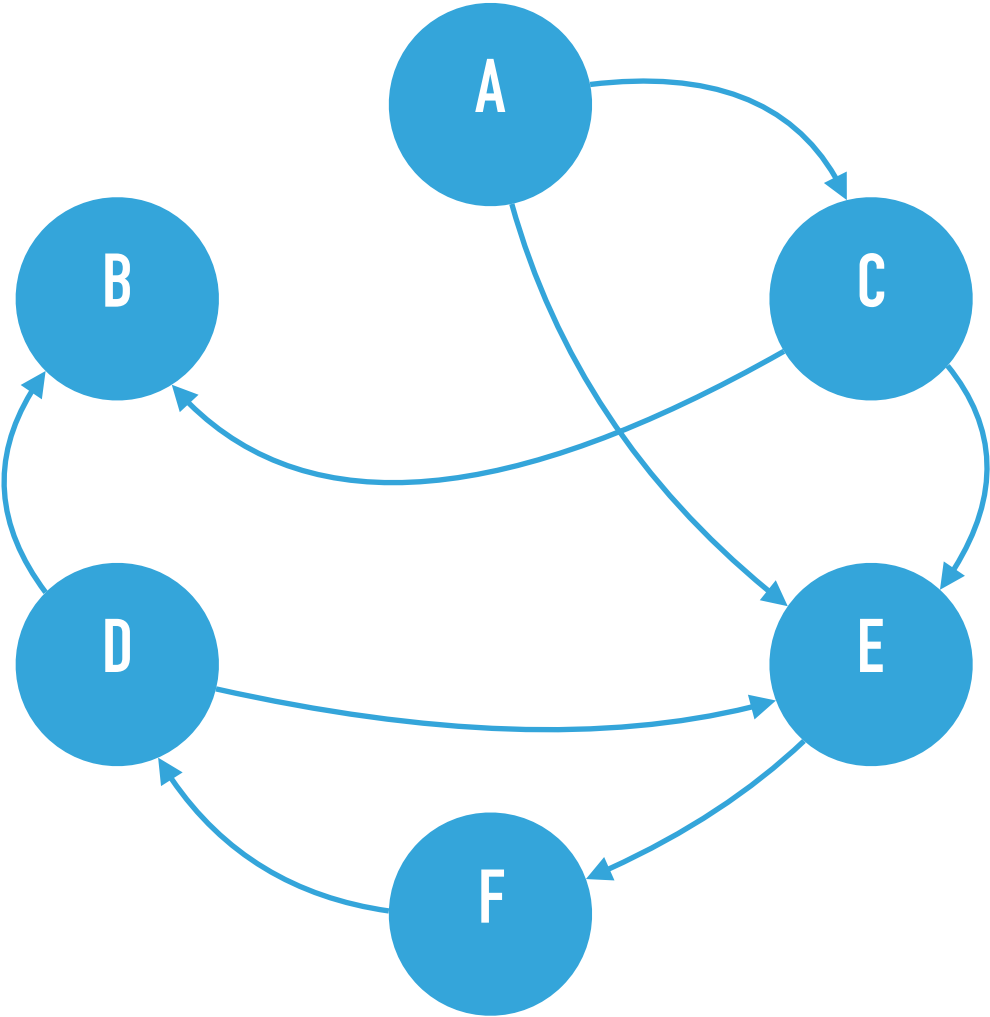


BREADTH FIRST SEARCH

- ▶ Find a path from A to F using breadth first search
 - ▶ (Assume that nodes are pushed onto the queue in *alphabetic order*)
- ▶ $A \rightarrow C \rightarrow B$
- ▶ Is *this* the shortest path?
 - ▶ Yes



BREADTH FIRST SEARCH

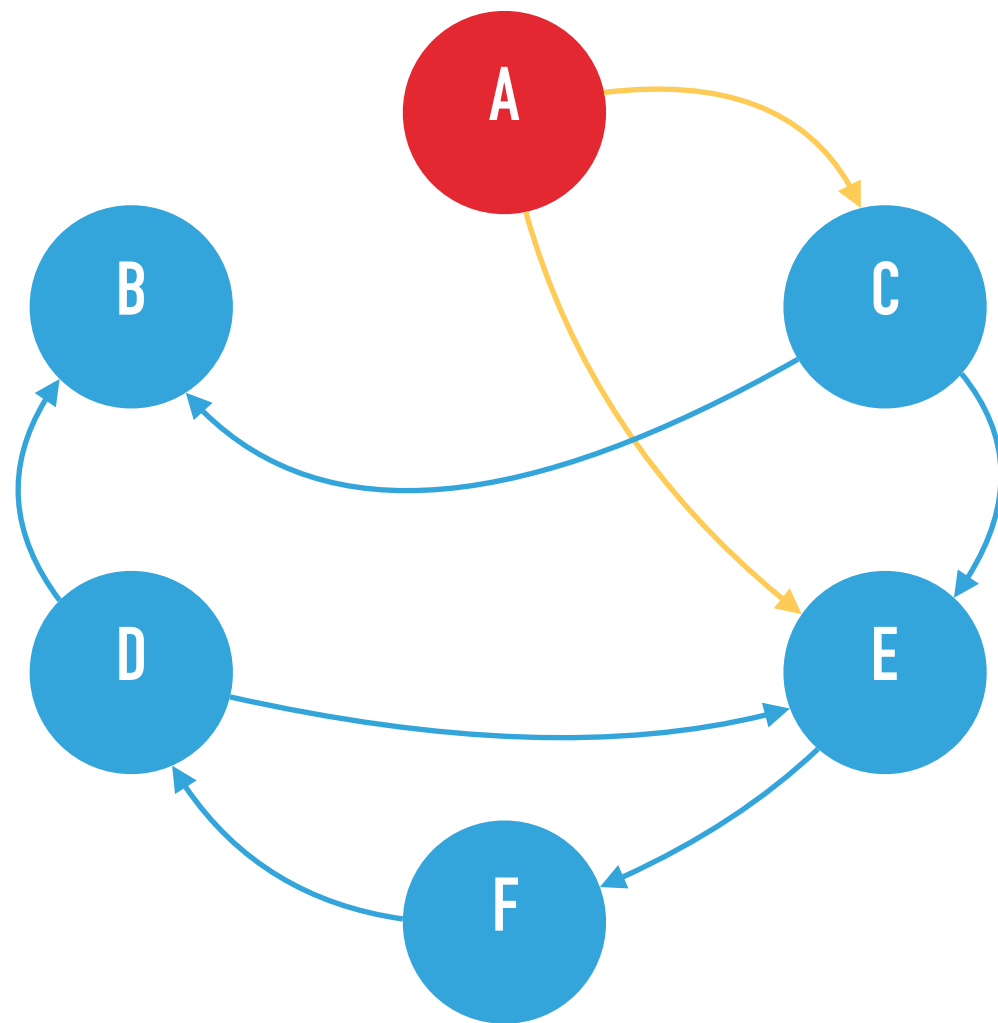


Paths to Consider (Queue)

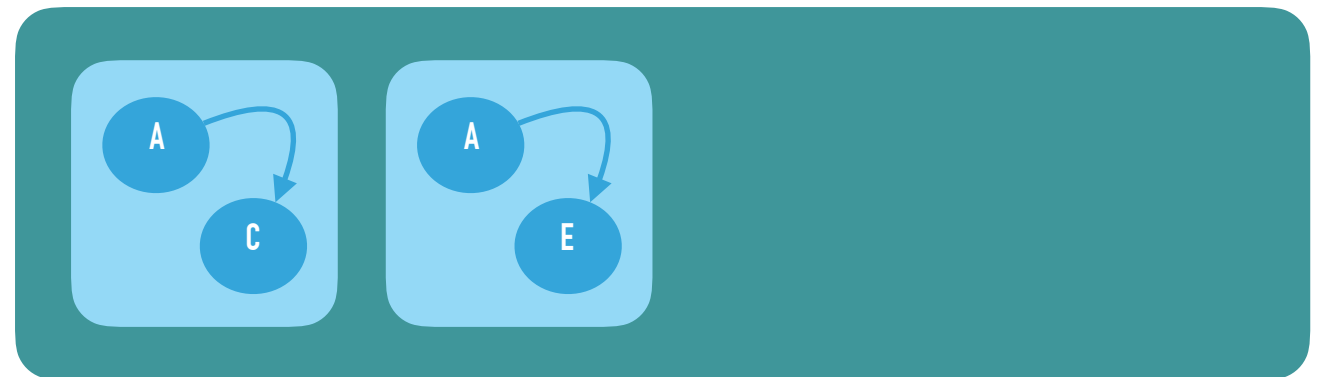


Current Path

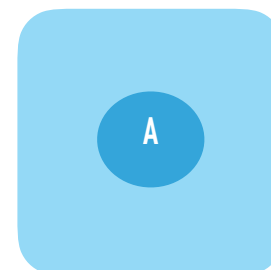
BREADTH FIRST SEARCH



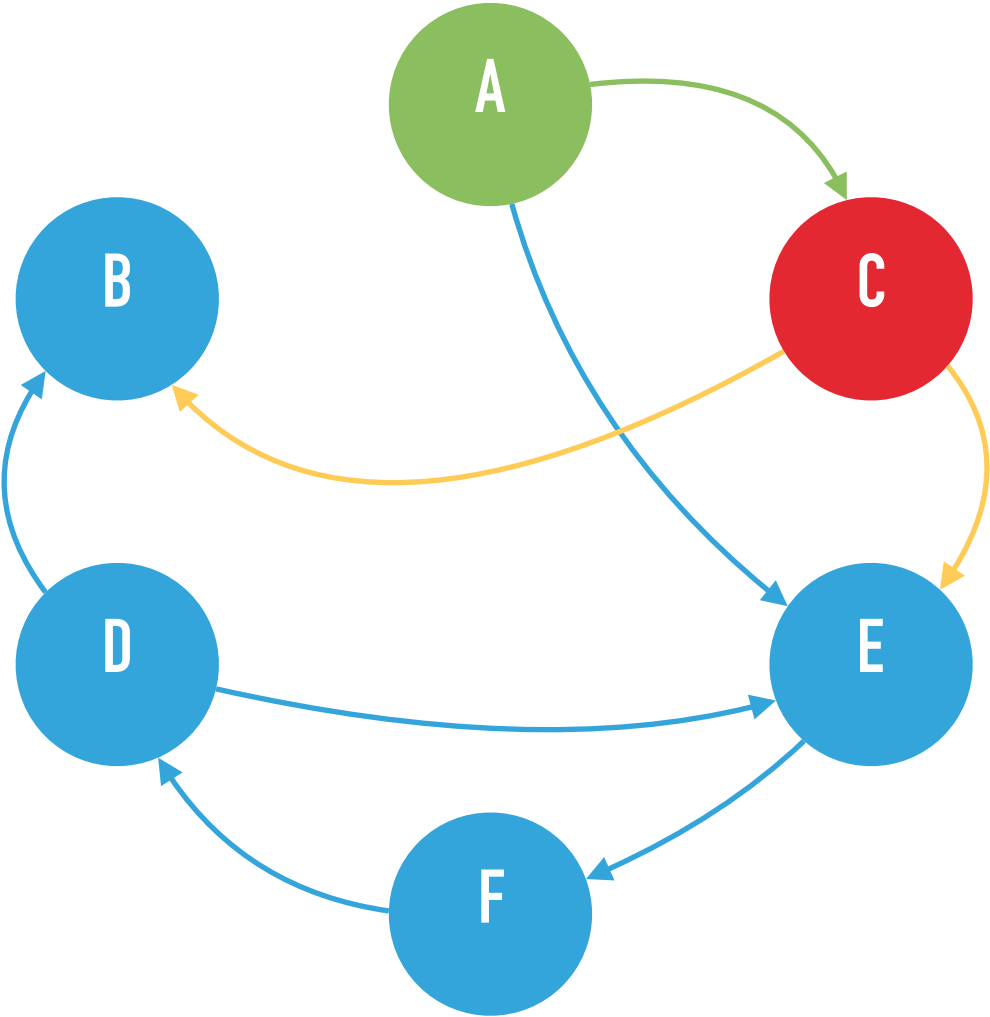
Paths to Consider (Queue)



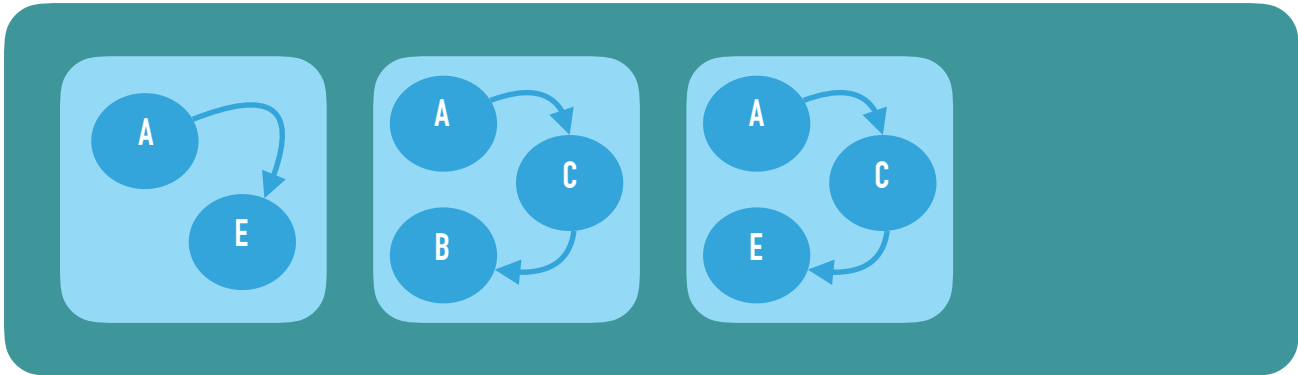
Current Path



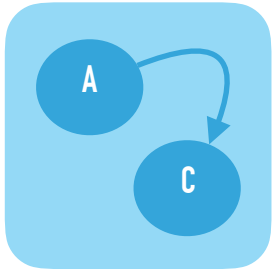
BREADTH FIRST SEARCH



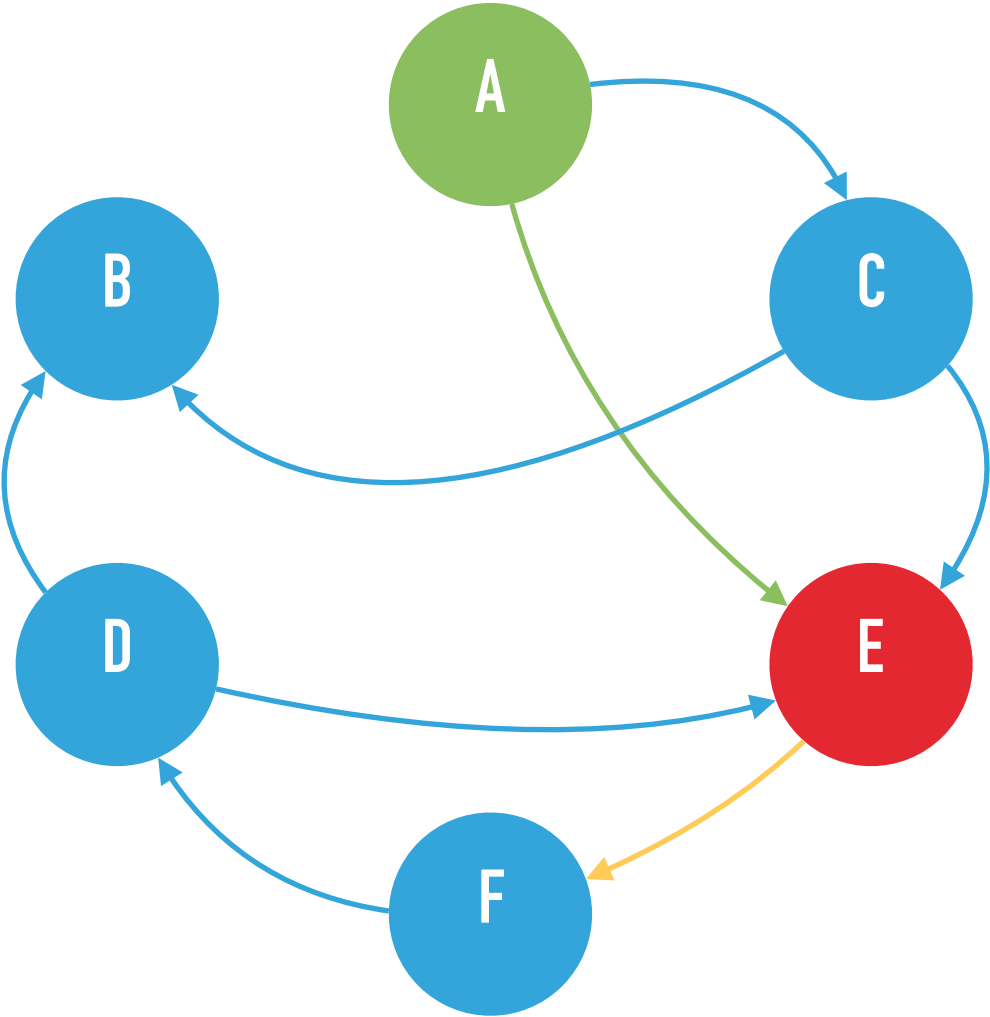
Paths to Consider (Queue)



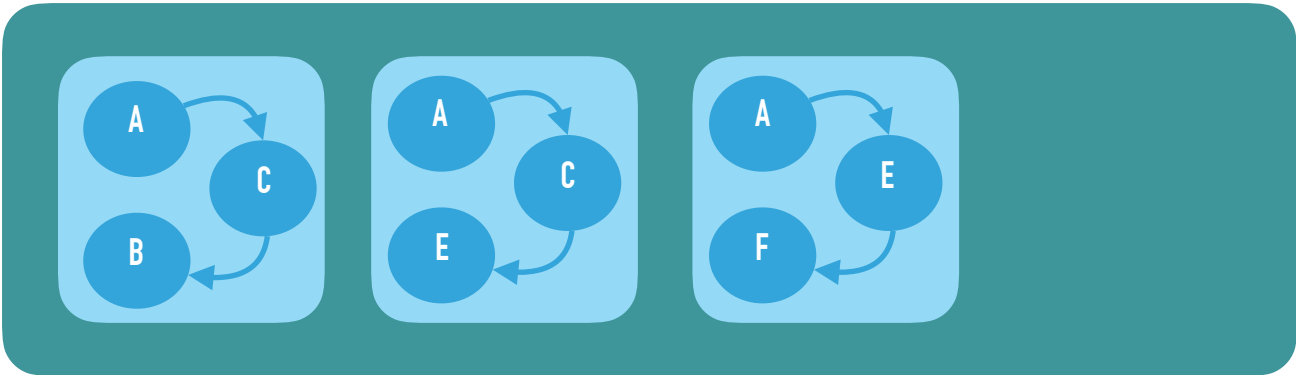
Current Path



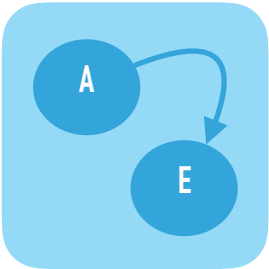
BREADTH FIRST SEARCH



Paths to Consider (Queue)



Current Path



**YOU NEVER CONSIDER A PATH OF
LENGTH $K + 1$**

**UNTIL YOU'VE CONSIDERED ALL PATHS OF
LENGTH K OR SHORTER**

COMPARING DFS AND BFS

COMPARING DFS AND BFS

DFS

- ▶ create a path with just start node and push onto stack *s*
- ▶ while *s* is not empty:
 - ▶ *p* = *s*.pop()
 - ▶ *v* = last node of *p*
 - ▶ if *v* is end node, you're done
 - ▶ mark *v* as visited
 - ▶ for each unvisited neighbor:
 - ▶ create new path and append neighbor
 - ▶ push new path onto *s*

BFS

- ▶ create a path with just start node and enqueue into queue *q*
- ▶ while *q* is not empty:
 - ▶ *p* = *q*.dequeue()
 - ▶ *v* = last node of *p*
 - ▶ if *v* is end node, you're done
 - ▶ mark *v* as visited
 - ▶ for each unvisited neighbor:
 - ▶ create new path and append neighbor
 - ▶ enqueue new path into *q*

COMPARING DFS AND BFS

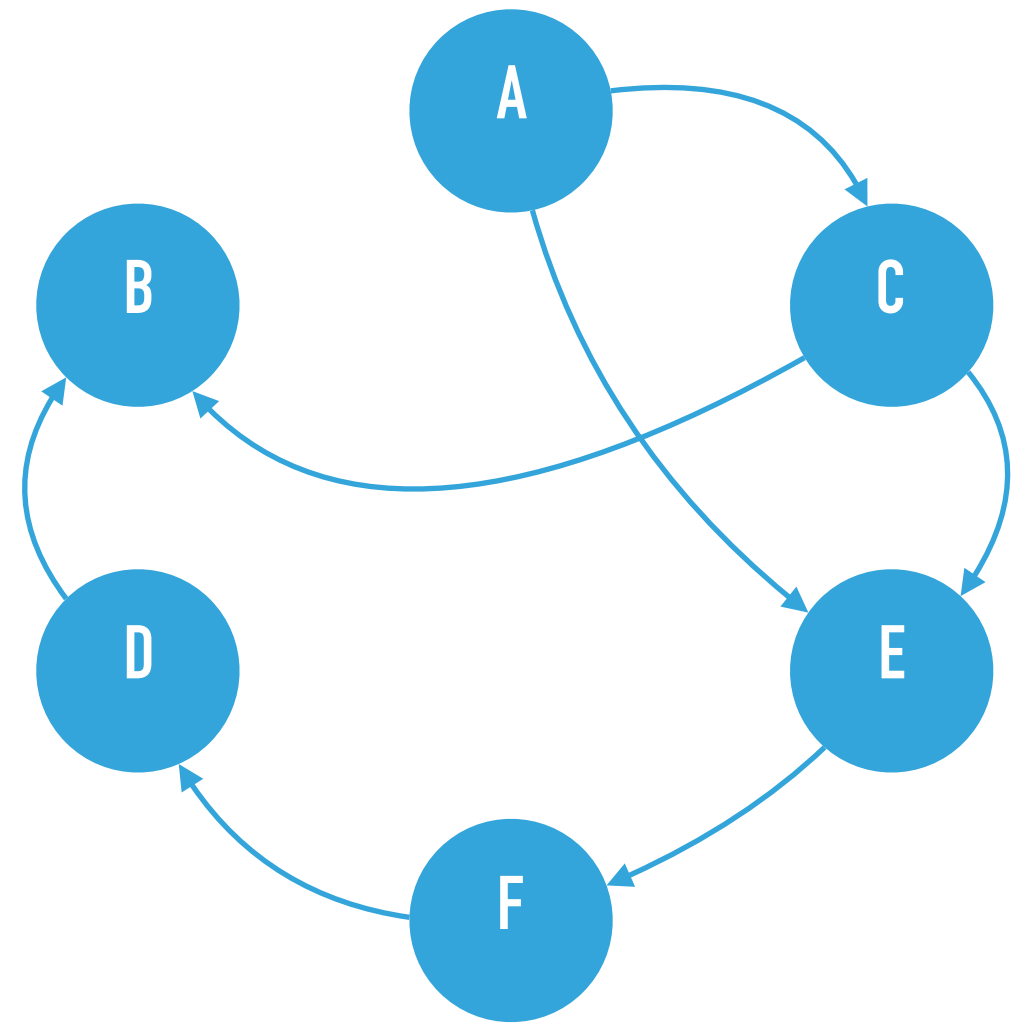
DFS

- ▶ create a path with just start node and push onto **stack s**
- ▶ while **s** is not empty:
 - ▶ **p = s.pop()**
 - ▶ **v = last node of p**
 - ▶ if **v** is end node, you're done
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 - ▶ **push new path onto s**

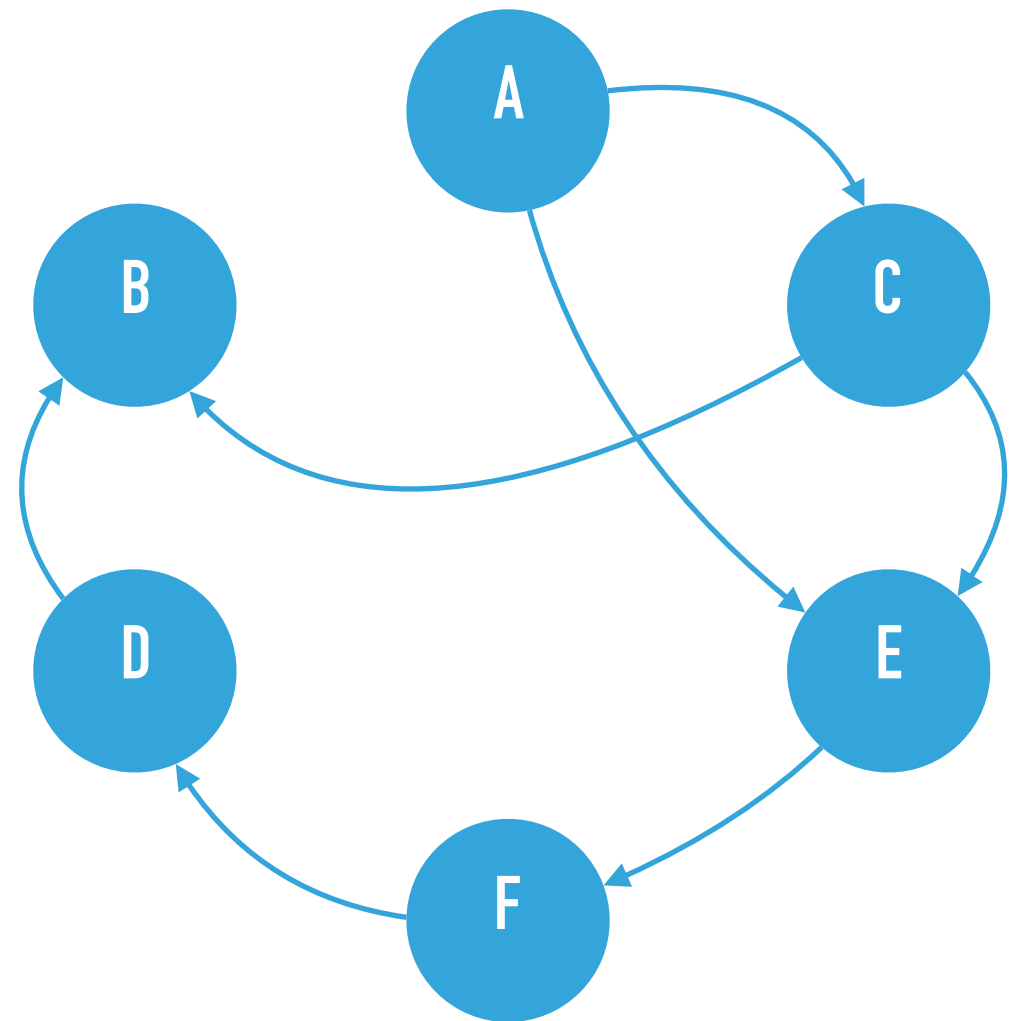
BFS

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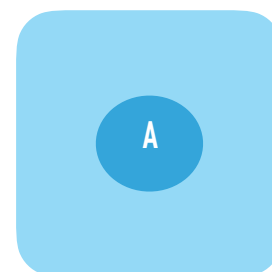
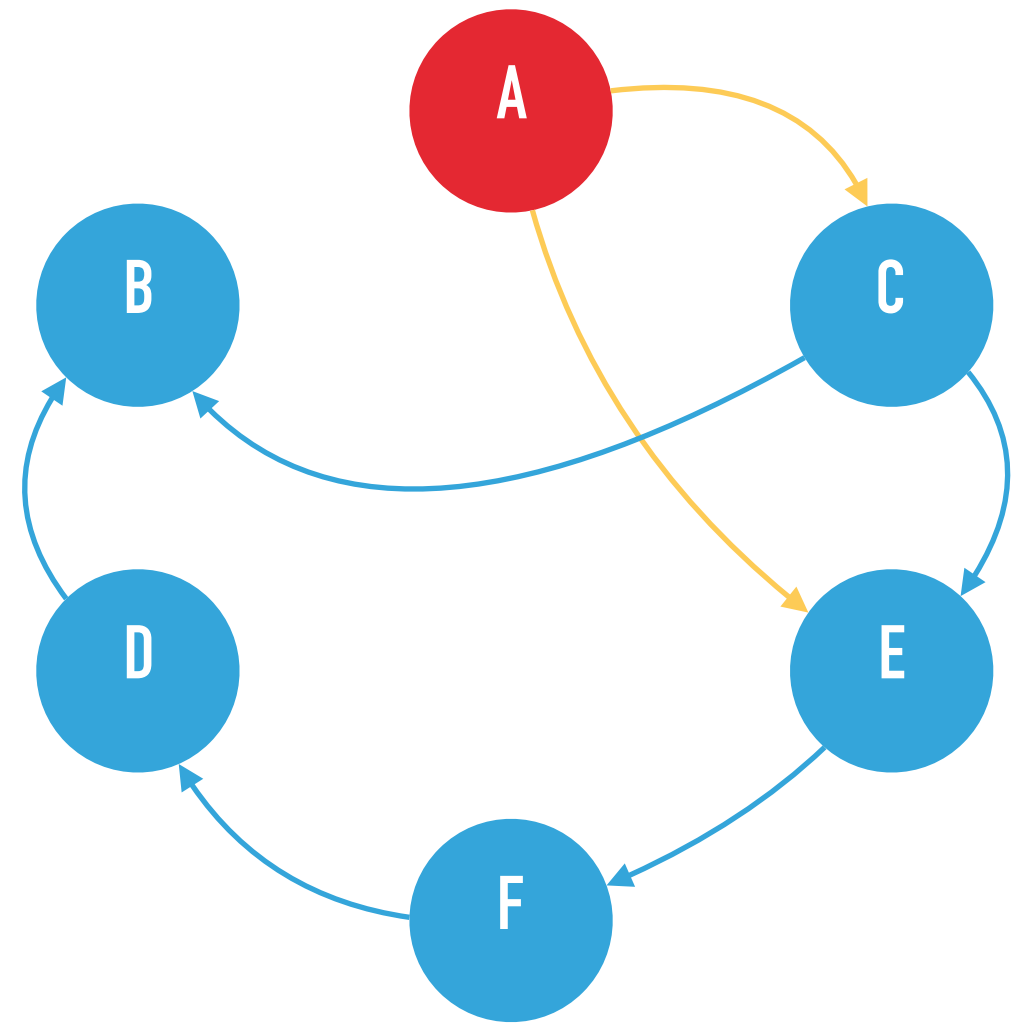
THE GRAPH SEARCH TO-DO LIST



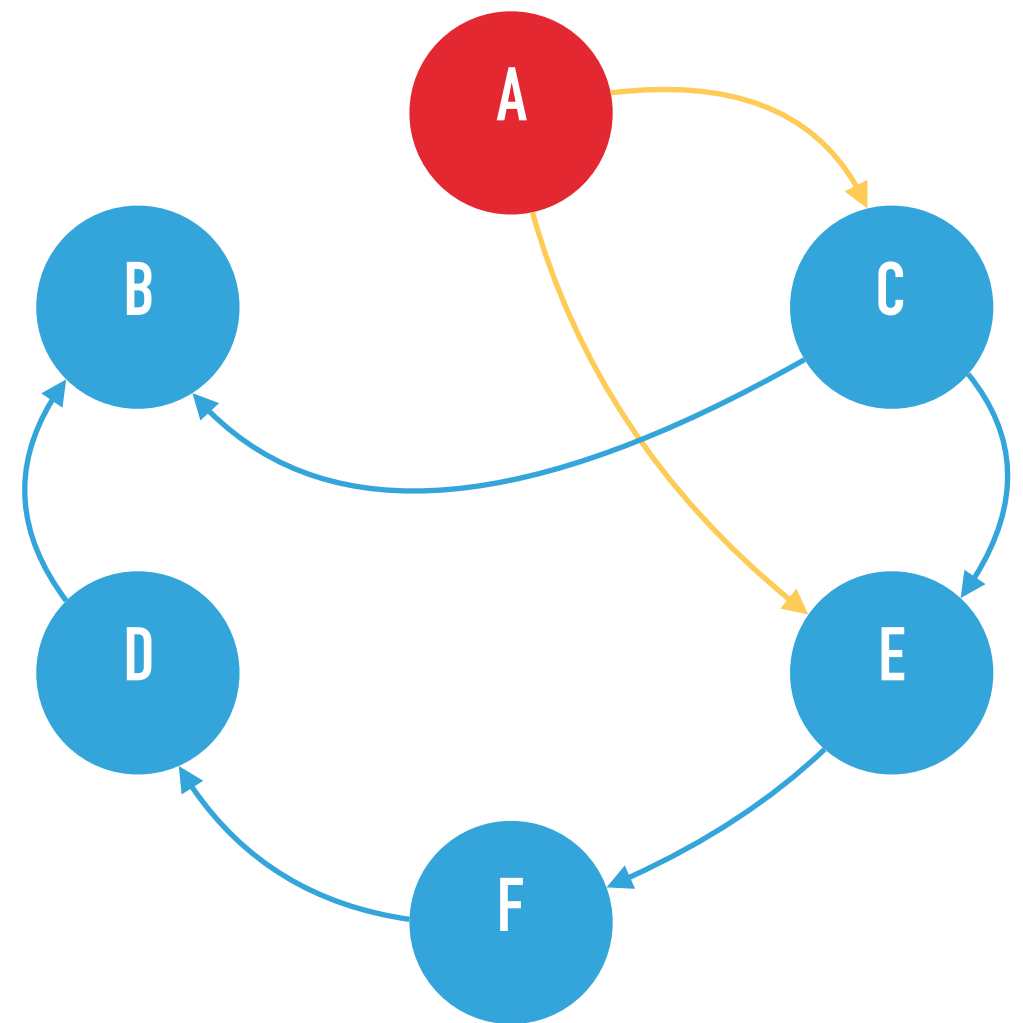
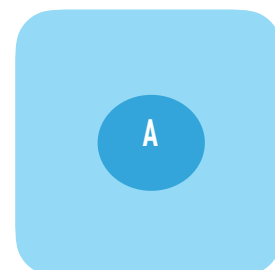
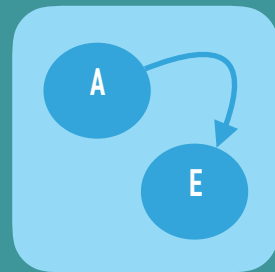
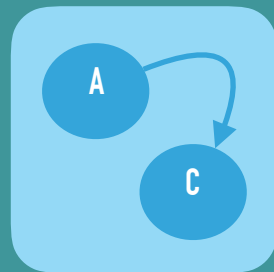
THE GRAPH SEARCH TO-DO LIST



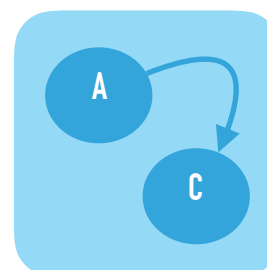
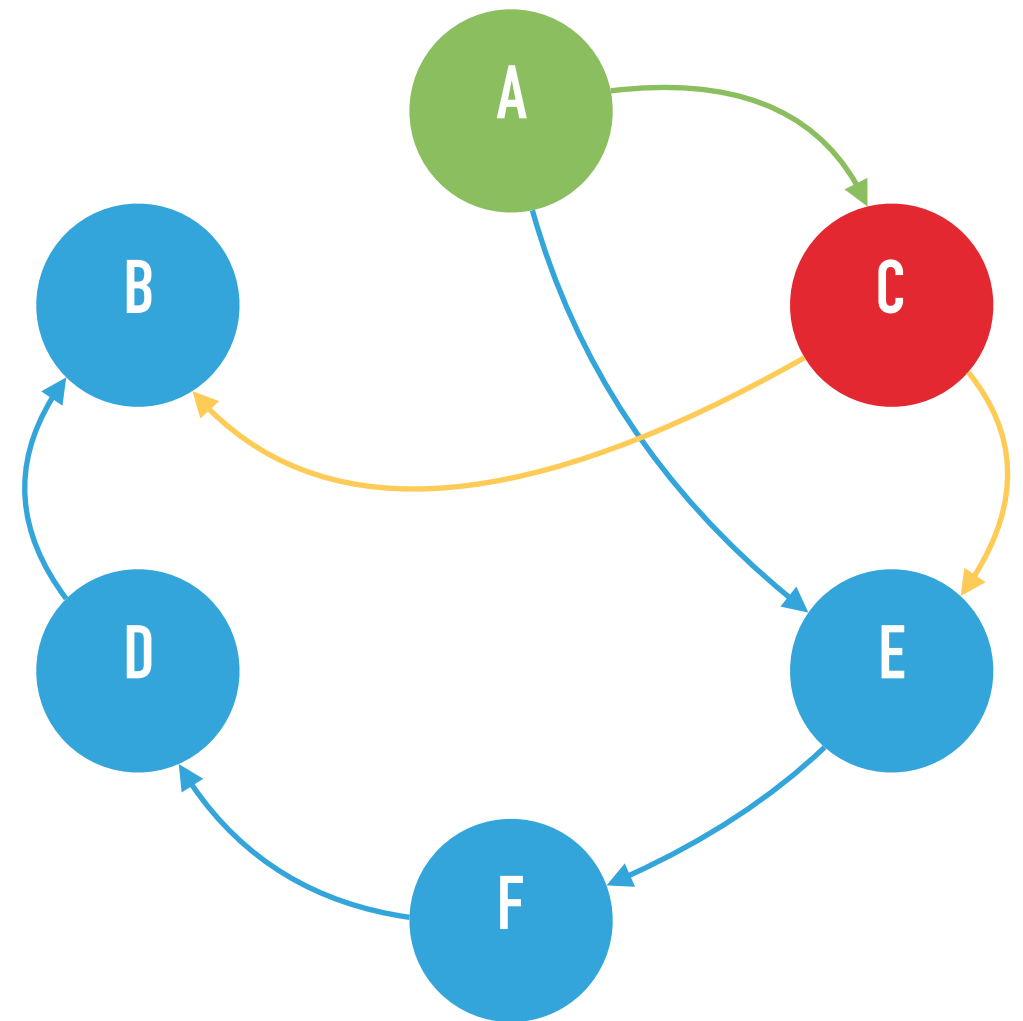
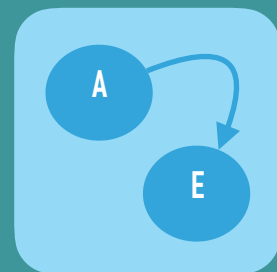
THE GRAPH SEARCH TO-DO LIST



THE GRAPH SEARCH TO-DO LIST

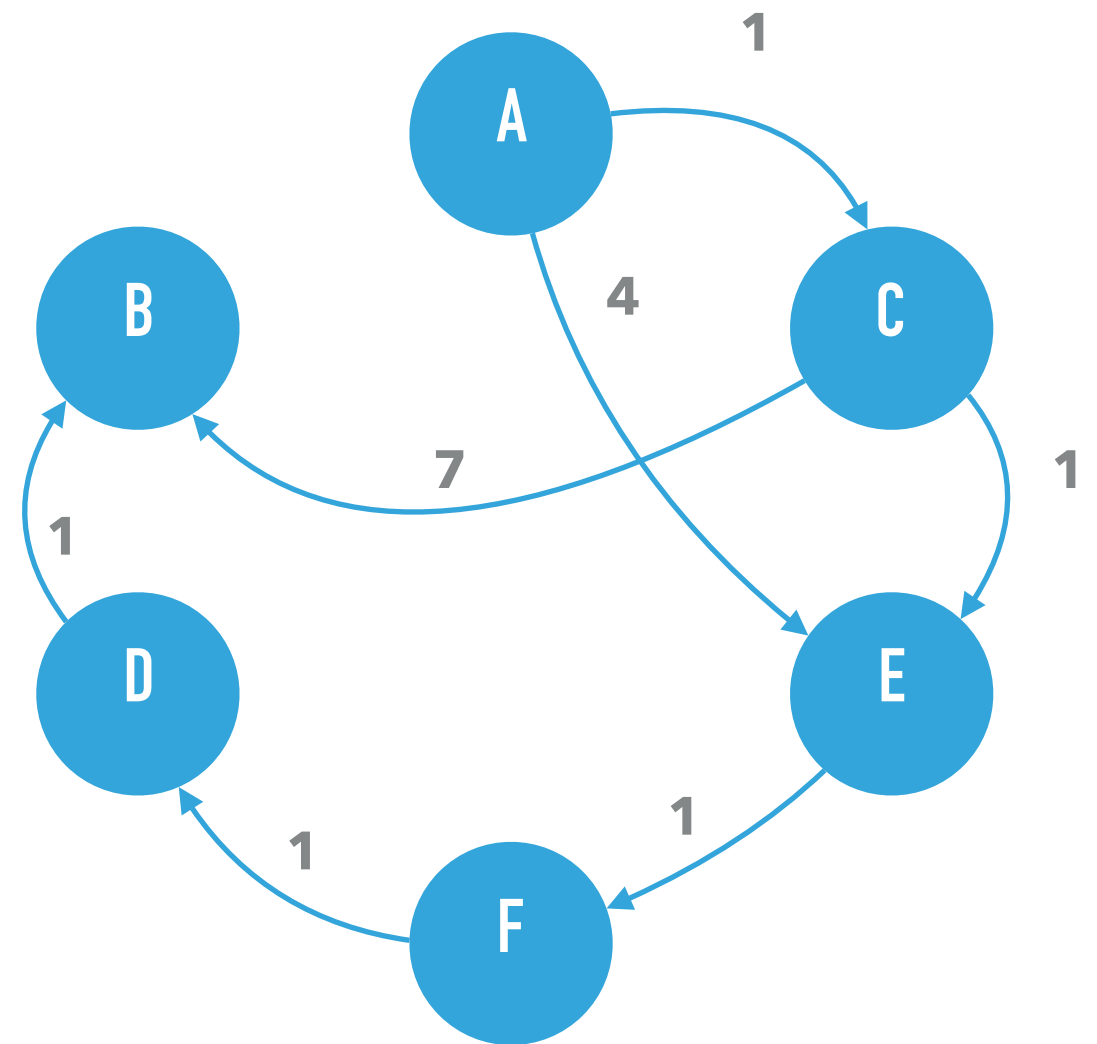


THE GRAPH SEARCH TO-DO LIST

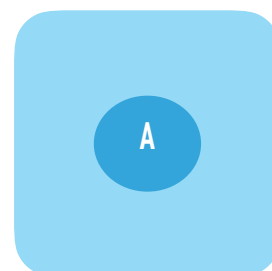
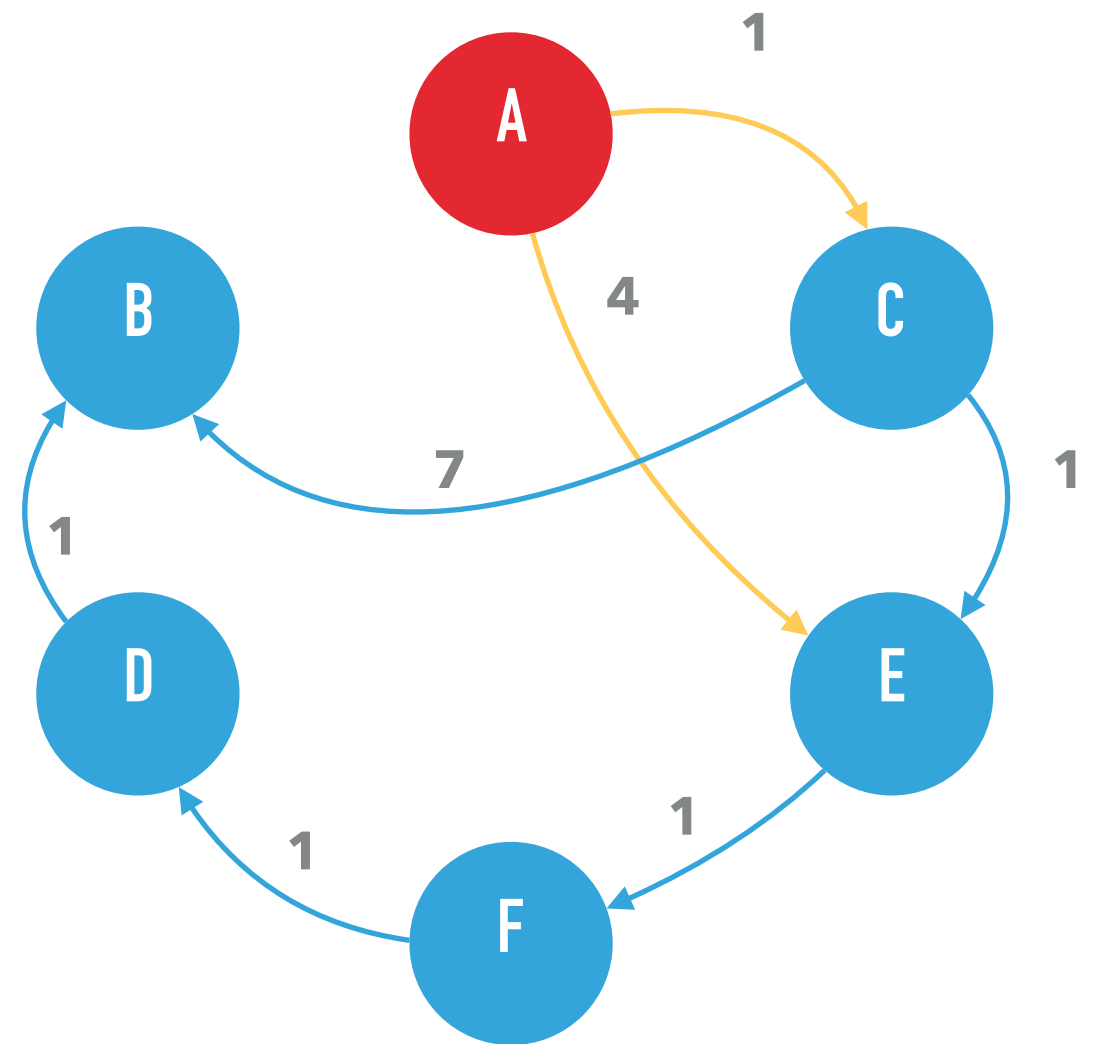


WEIGHTY DECISIONS

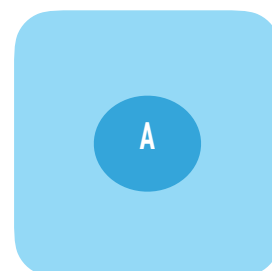
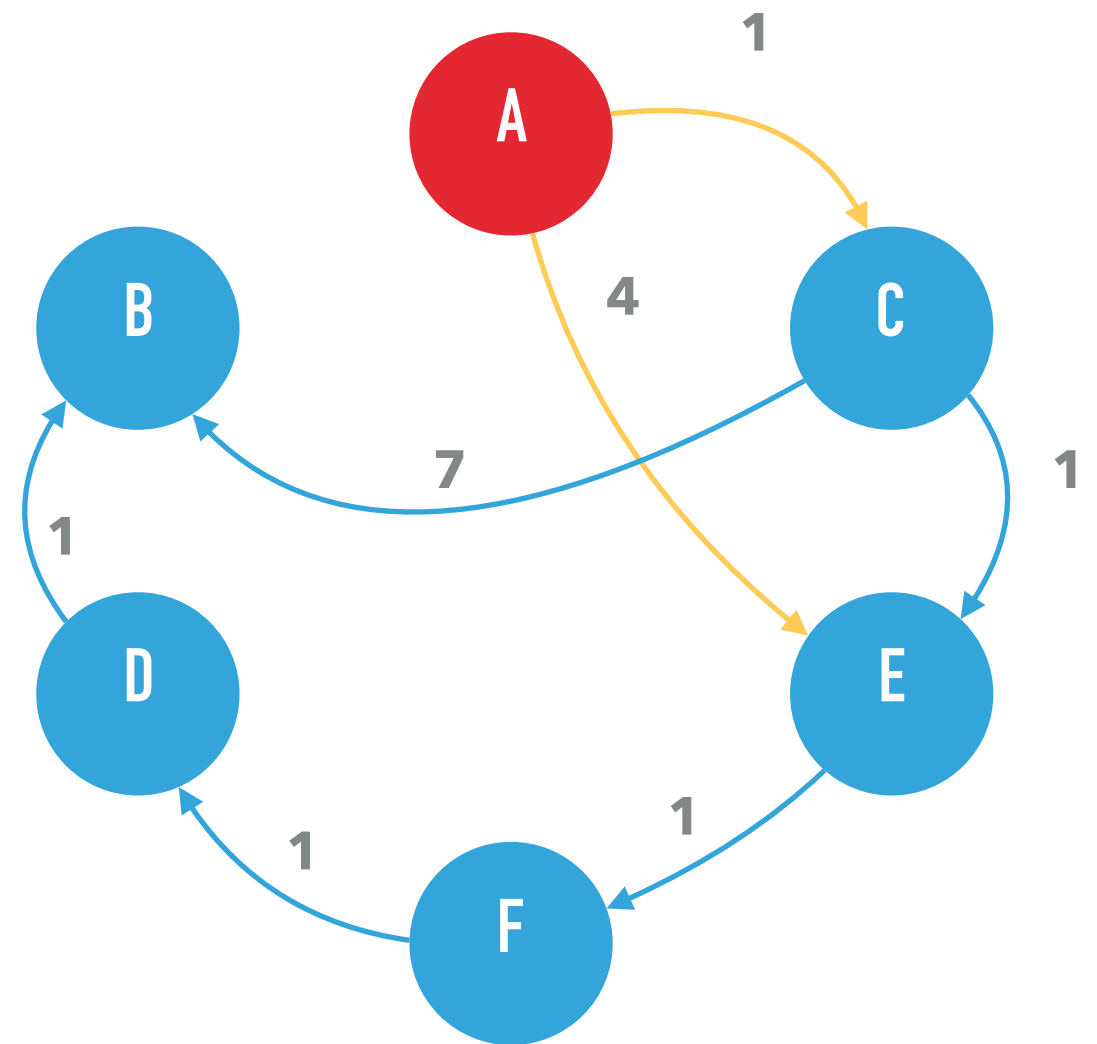
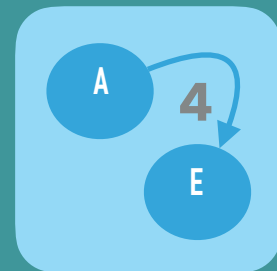
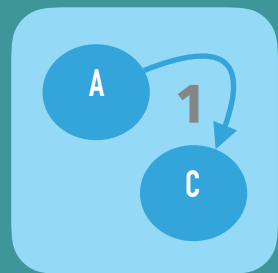
DEALING WITH WEIGHTY TOPICS



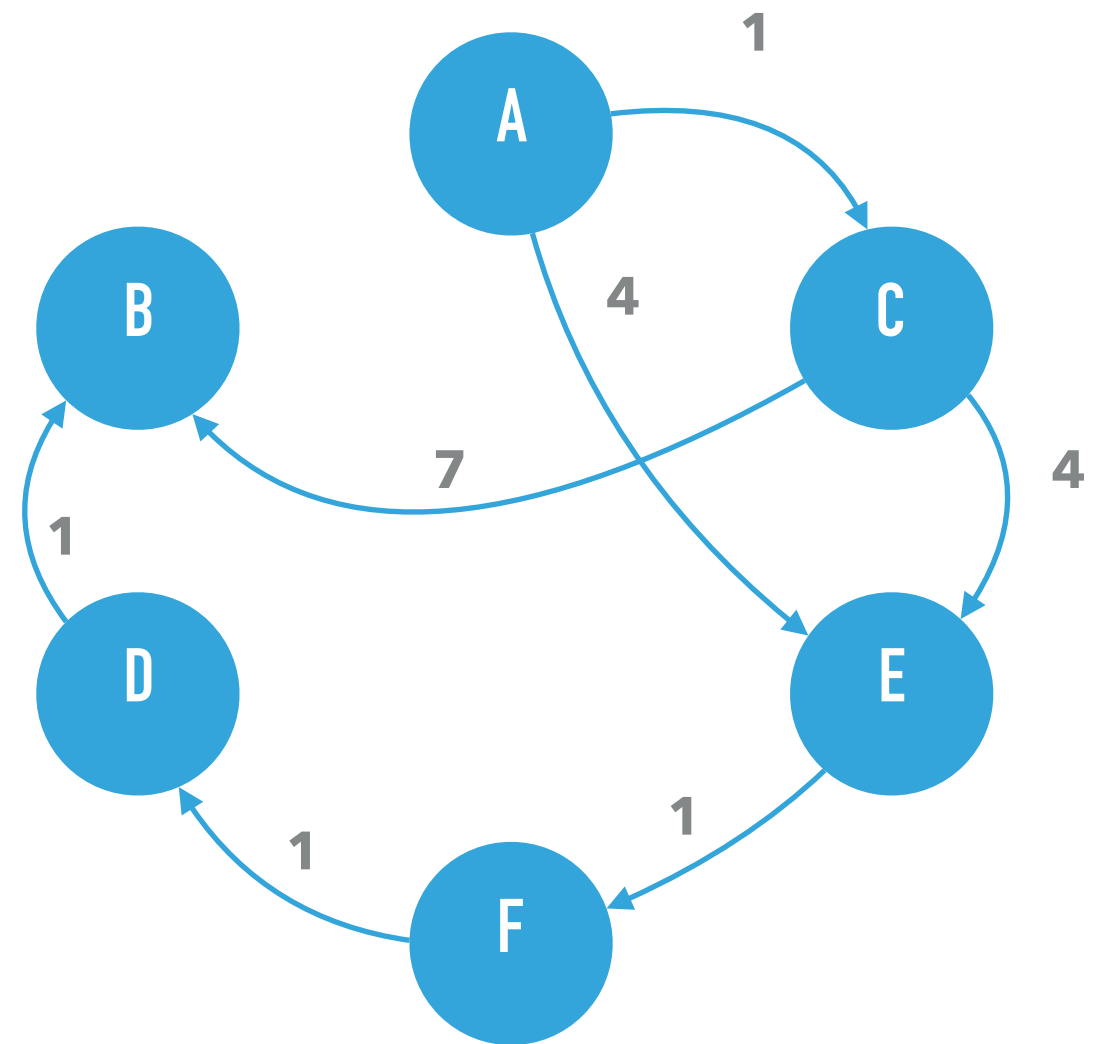
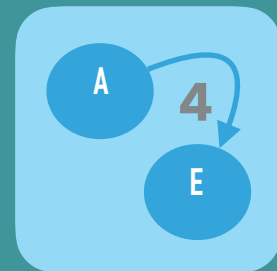
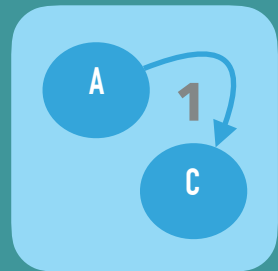
DEALING WITH WEIGHTY TOPICS



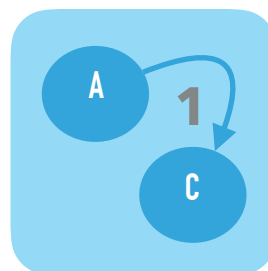
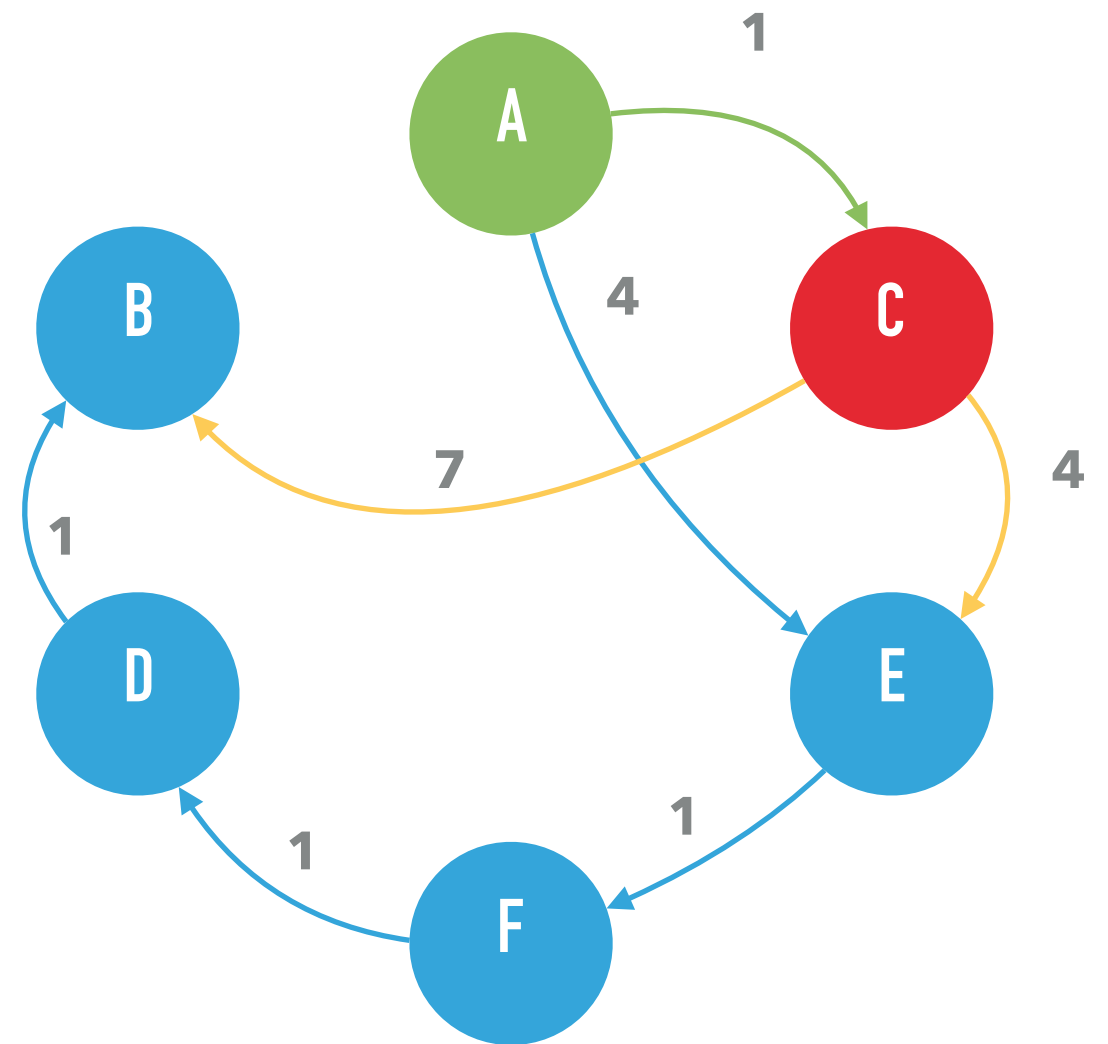
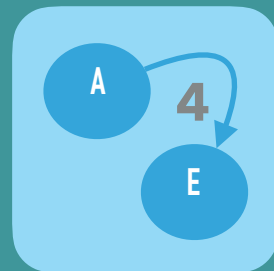
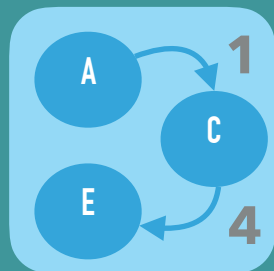
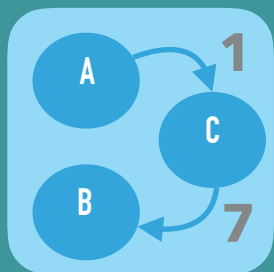
DEALING WITH WEIGHTY TOPICS



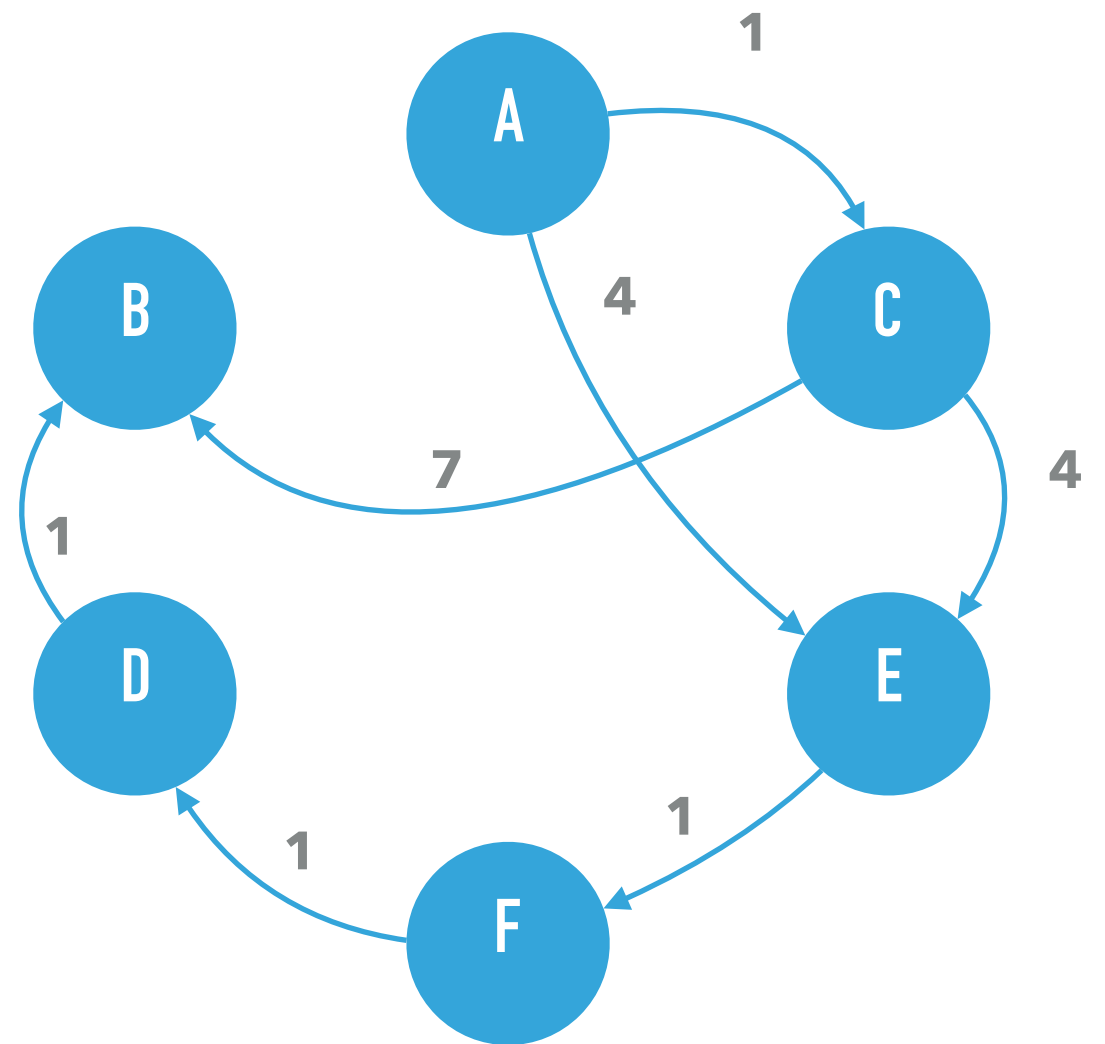
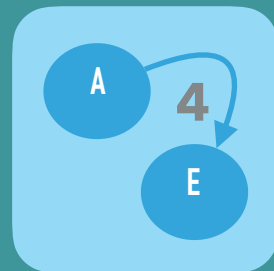
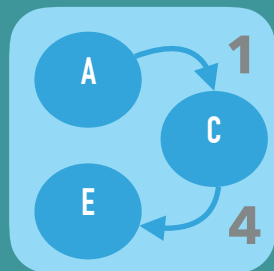
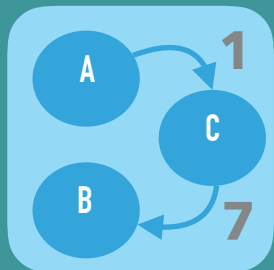
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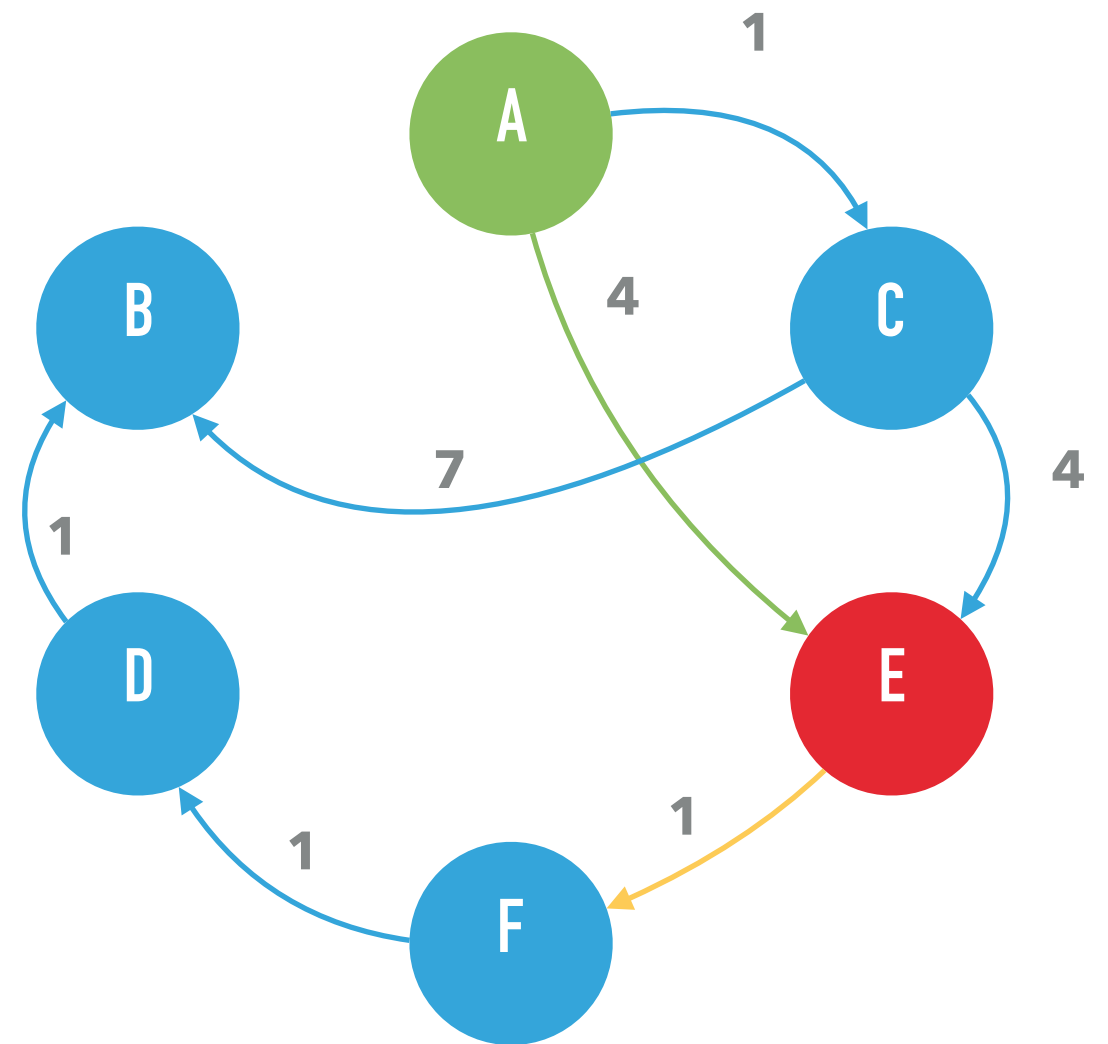
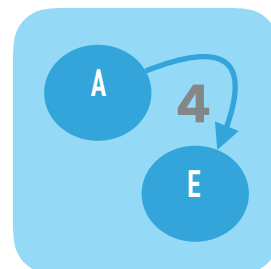
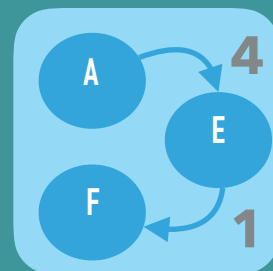
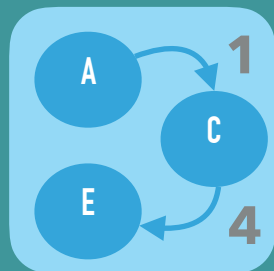
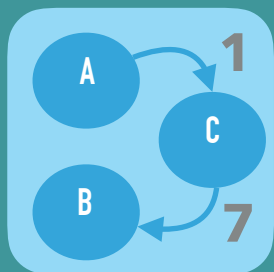
DEALING WITH WEIGHTY TOPICS



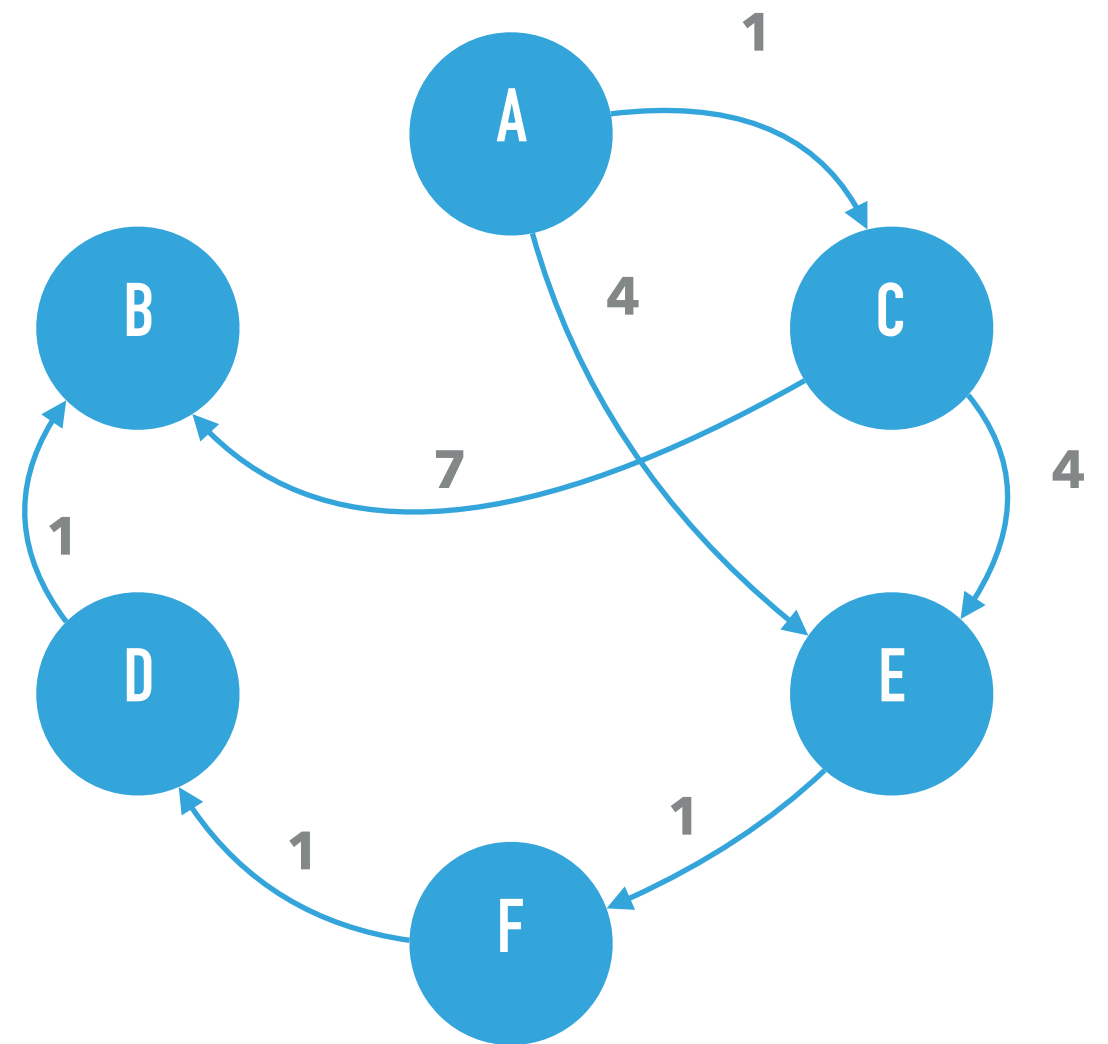
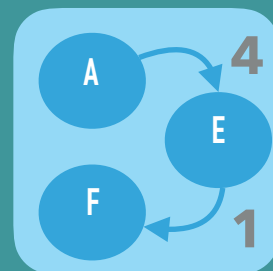
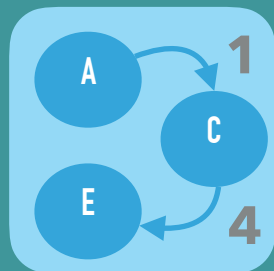
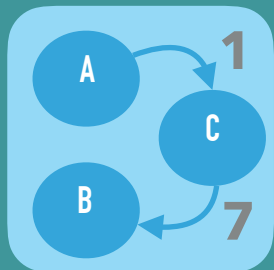
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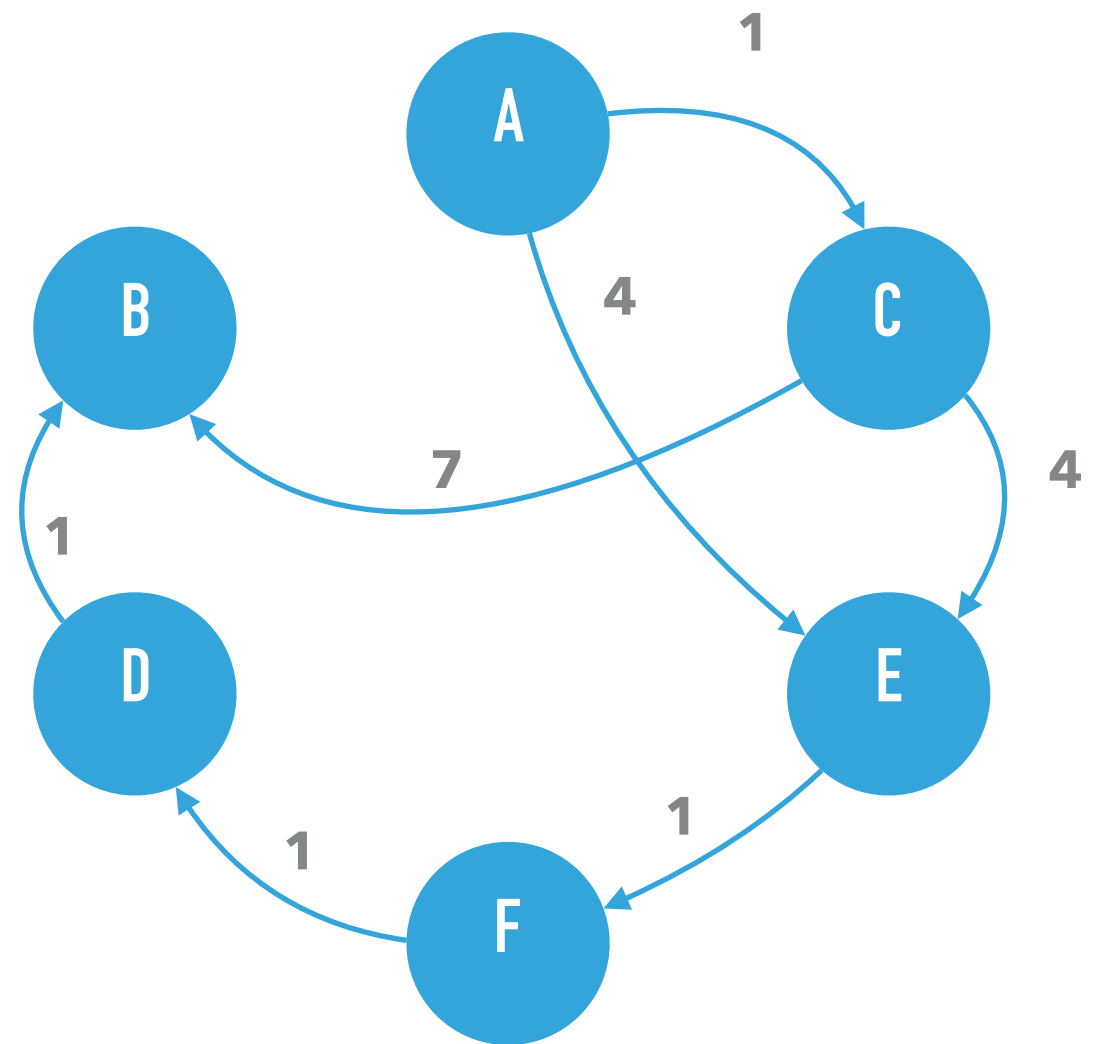
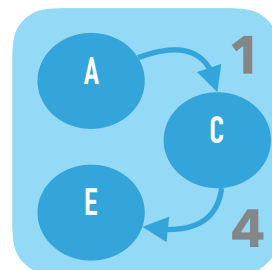
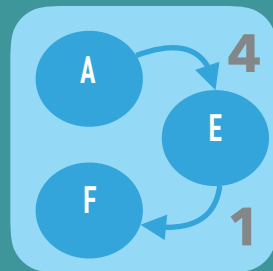
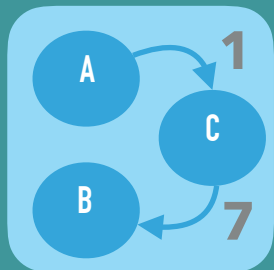
DEALING WITH WEIGHTY TOPICS



DEALING WITH WEIGHTY TOPICS



DEALING WITH WEIGHTY TOPICS



**IN DIJKSTRA'S ALGORITHM,
THE TODO LIST IS A PRIORITY
QUEUE**

DIJKSTRA'S ALGORITHM (PSEUDOCODE)

- ▶ create a path with just start node and enqueue into priority queue q
- ▶ while q is not empty
 - ▶ $p = q.dequeue()$
 - ▶ $v = \text{last node of } p$
 - ▶ if v is end node, you're done
 - ▶ if you've seen v before, skip it
 - ▶ mark v as visited
 - ▶ for each unvisited neighbor:
 - ▶ create new path and append neighbor
 - ▶ enqueue new path into q

DIJKSTRA'S ALGORITHM (PSEUDOCODE)

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 - ▶ for each unvisited neighbor:
 - ▶ create new path and append neighbor
 - ▶ enqueue new path into q with priority pathLength

DIJKSTRA'S ODDS AND ENDS

- ▶ **create a path with just start node and enqueue into priority queue q**
- ▶ while q is not empty
 - ▶ $p = q.dequeue()$
 - ▶ $v = \text{last node of } p$
 - ▶ if v is end node, you're done
 - ▶ if you've seen v before, skip it
 - ▶ mark v as visited
 - ▶ for each unvisited neighbor:
 - ▶ create new path and append neighbor
 - ▶ enqueue new path into q with priority `pathLength`
- ▶ What do you initialize the weight of the path to?

DIJKSTRA'S ODDS AND ENDS

- ▶ **create a path with just start node and enqueue into priority queue q**
- ▶ while q is not empty
 - ▶ $p = q.dequeue()$
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 - ▶ mark v as visited
 - ▶ for each unvisited neighbor:
 - ▶ create new path and append neighbor
 - ▶ enqueue new path into q with priority `pathLength`
- ▶ What do you initialize the weight of the path to?
 - ▶ Zero should be fine

DIJKSTRA'S ODDS AND ENDS

- ▶ create a path with just start node and enqueue into priority queue q
- ▶ while q is not empty
 - ▶ $p = q.dequeue()$
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 - ▶ mark v as visited
 - ▶ for each unvisited neighbor:
 - ▶ create new path and append neighbor
 - ▶ enqueue new path into q with priority pathLength
- ▶ Can't I just return the path as soon as I find the end node? Why wait until I dequeue?

DIJKSTRA'S ODDS AND ENDS

- ▶ create a path with just start node and enqueue into priority queue q
- ▶ while q is not empty
 - ▶ `p = q.dequeue()`
 - ▶ `v = last node of p`
 - ▶ **if v is end node, you're done**
 - ▶ if you've seen v before, skip it
 - ▶ mark v as visited
 - ▶ for each unvisited neighbor:
 - ▶ create new path and append neighbor
 - ▶ enqueue new path into q with priority `pathLength`
- ▶ Can't I just return the path as soon as I find the end node? Why wait until I dequeue?
 - ▶ This is one of the most common mistakes people make with Dijkstra's!
 - ▶ It's possible a path with a lower priority gets enqueued in the meantime.

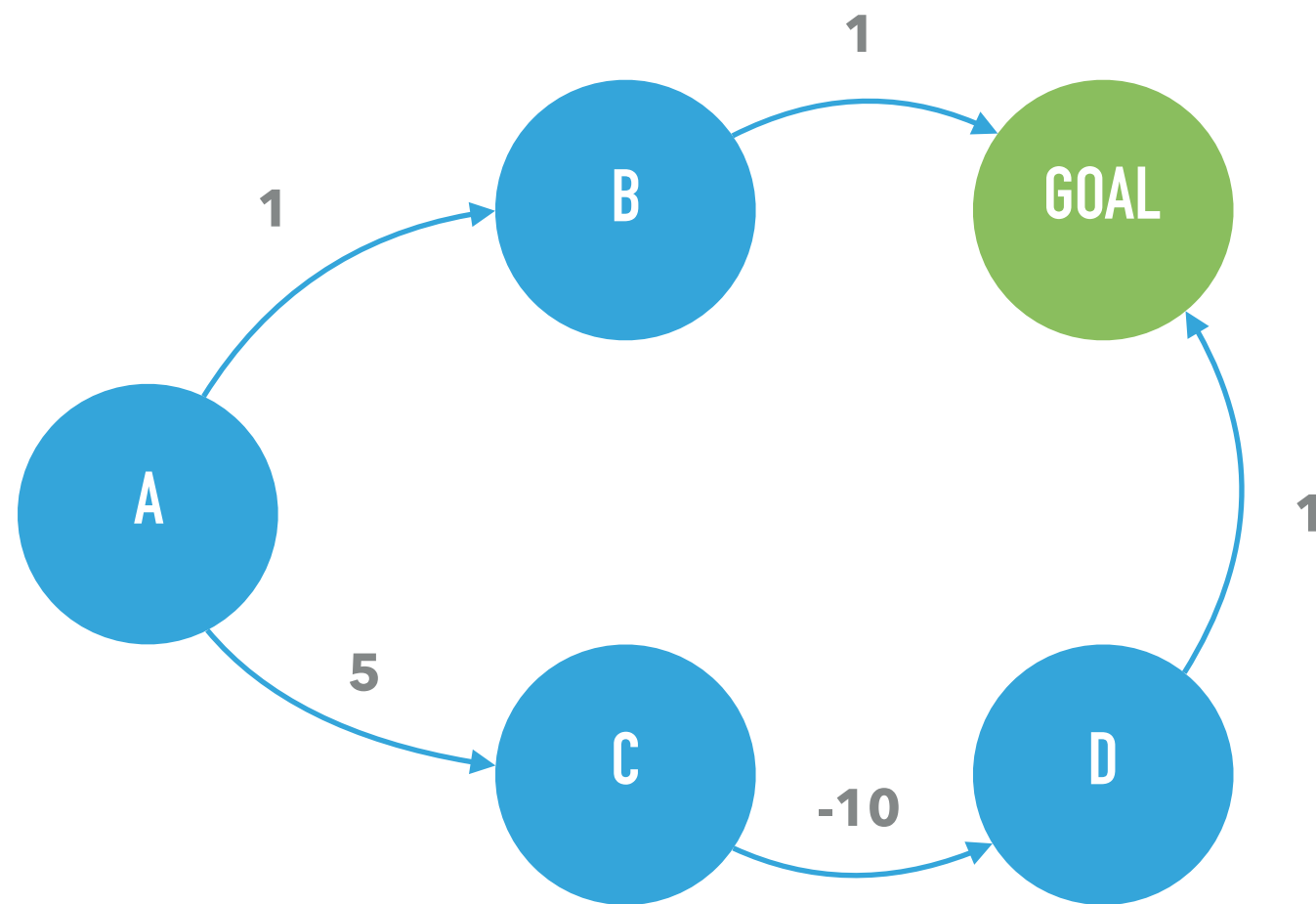
DIJKSTRA'S ODDS AND ENDS

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 - ▶ for each unvisited neighbor:
 - ▶ create new path and append neighbor
 - ▶ enqueue new path into q with priority pathLength
- ▶ Why would you skip the node just because you've seen it before?

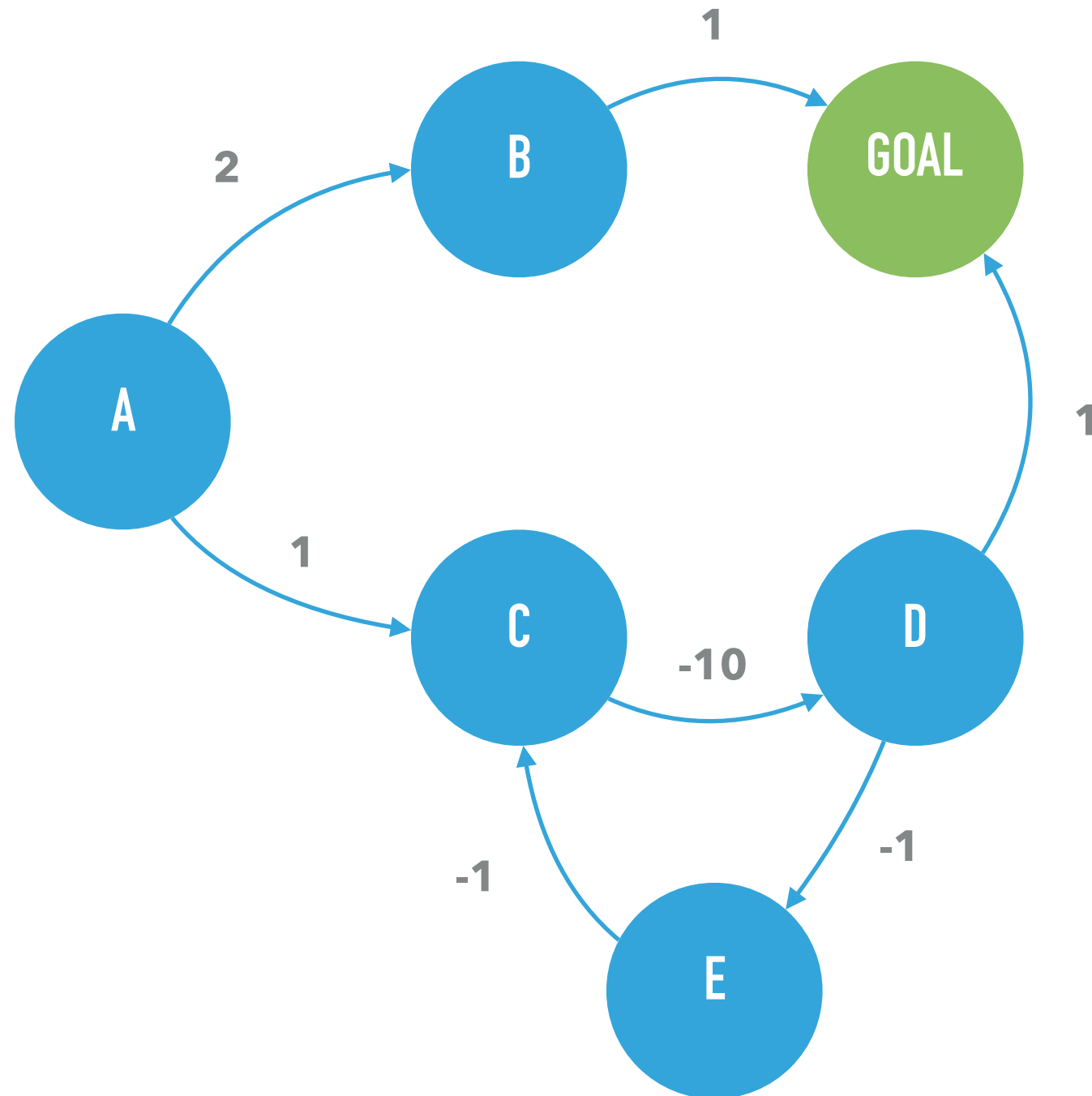
DIJKSTRA'S ODDS AND ENDS

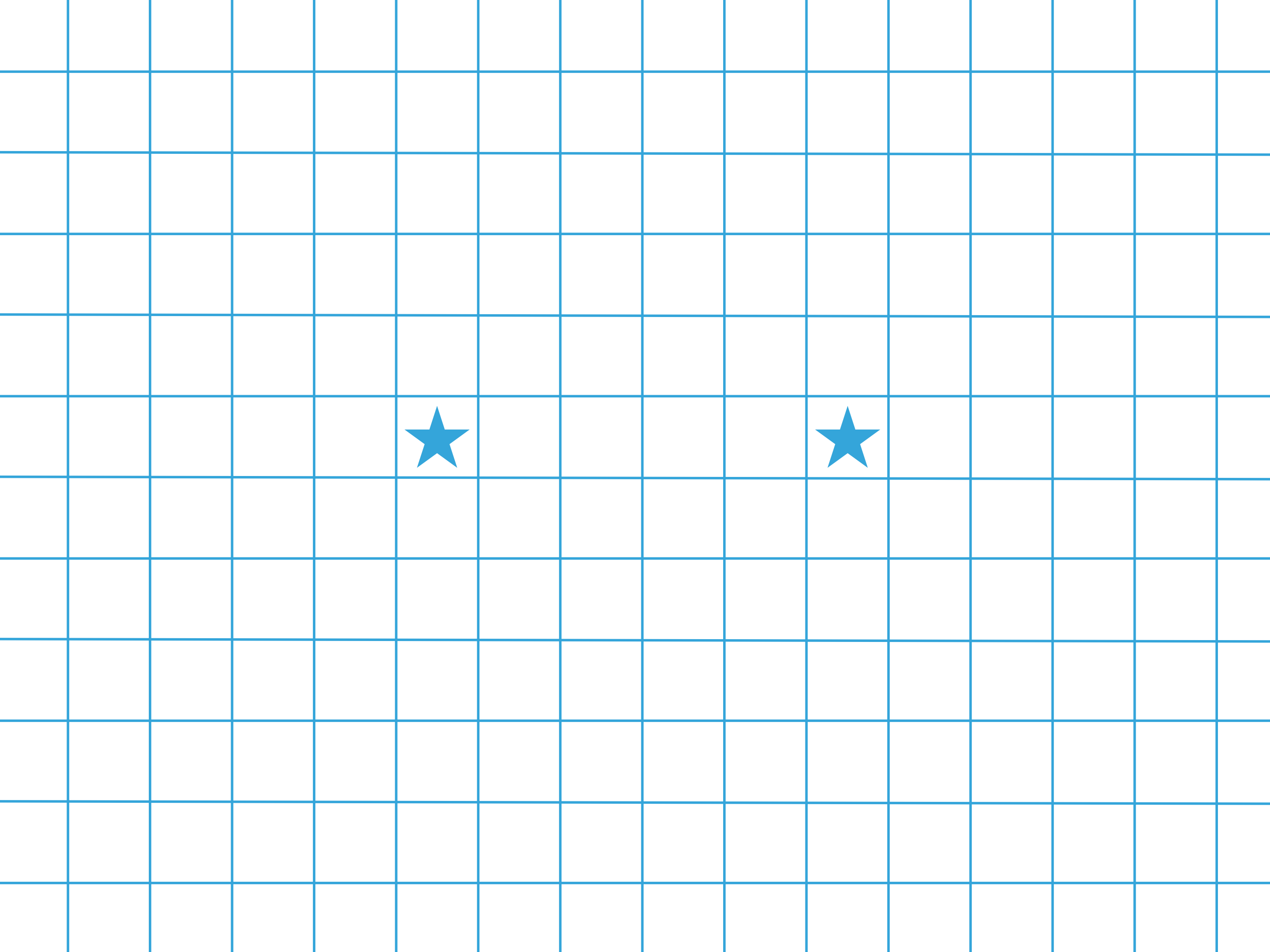
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 - ▶ mark v as visited
 - ▶ for each unvisited neighbor:
 - ▶ create new path and append neighbor
 - ▶ enqueue new path into q with priority pathLength
- ▶ Why would you skip the node just because you've seen it before?
 - ▶ If you've seen the node before, that means you've already found a shorter path to it.
 - ▶ Any path that follows from this one already has a shorter equivalent
 - ▶ **The first path you find to v will be the shortest path to v**

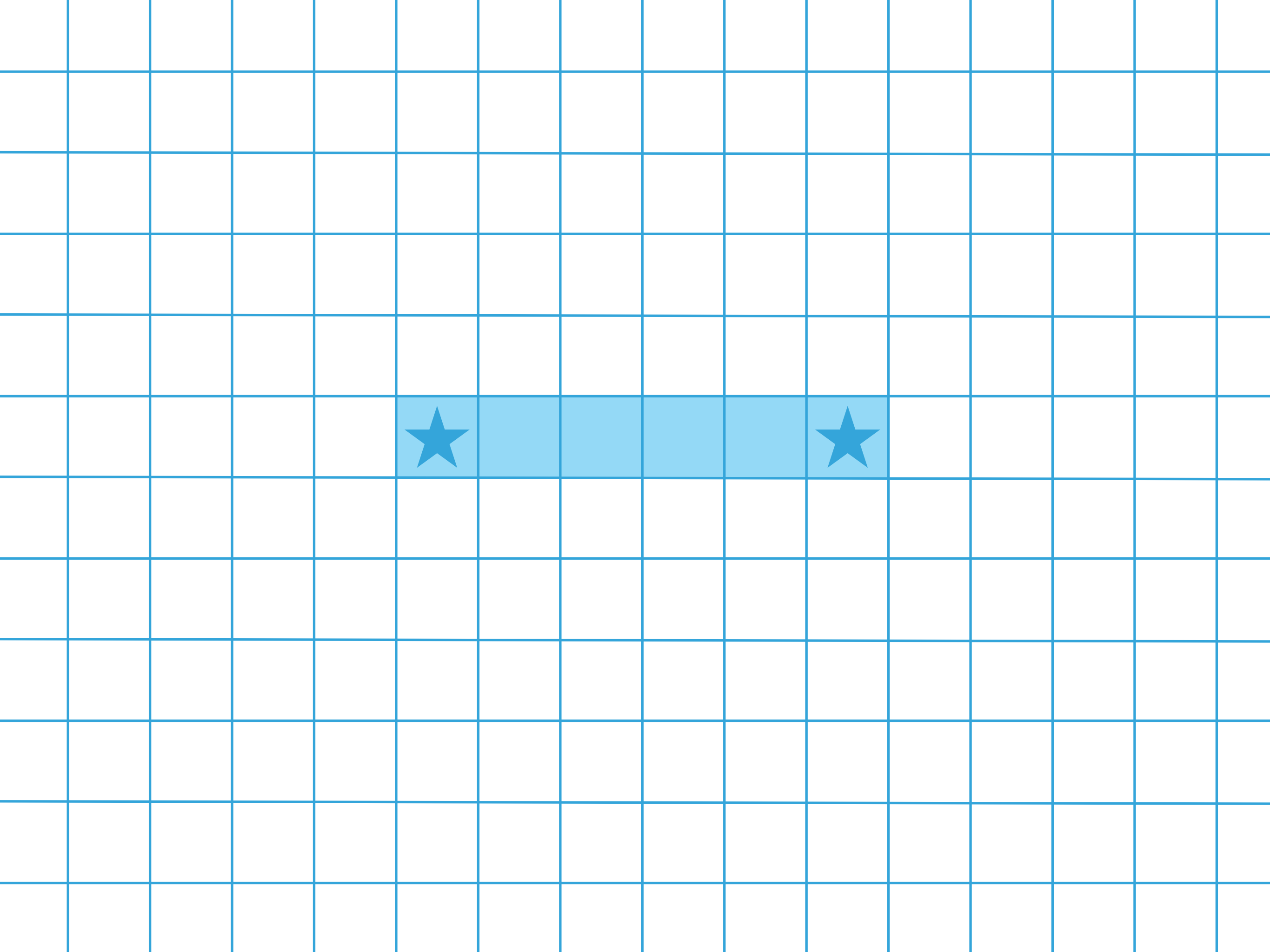
NEGATIVE EDGES

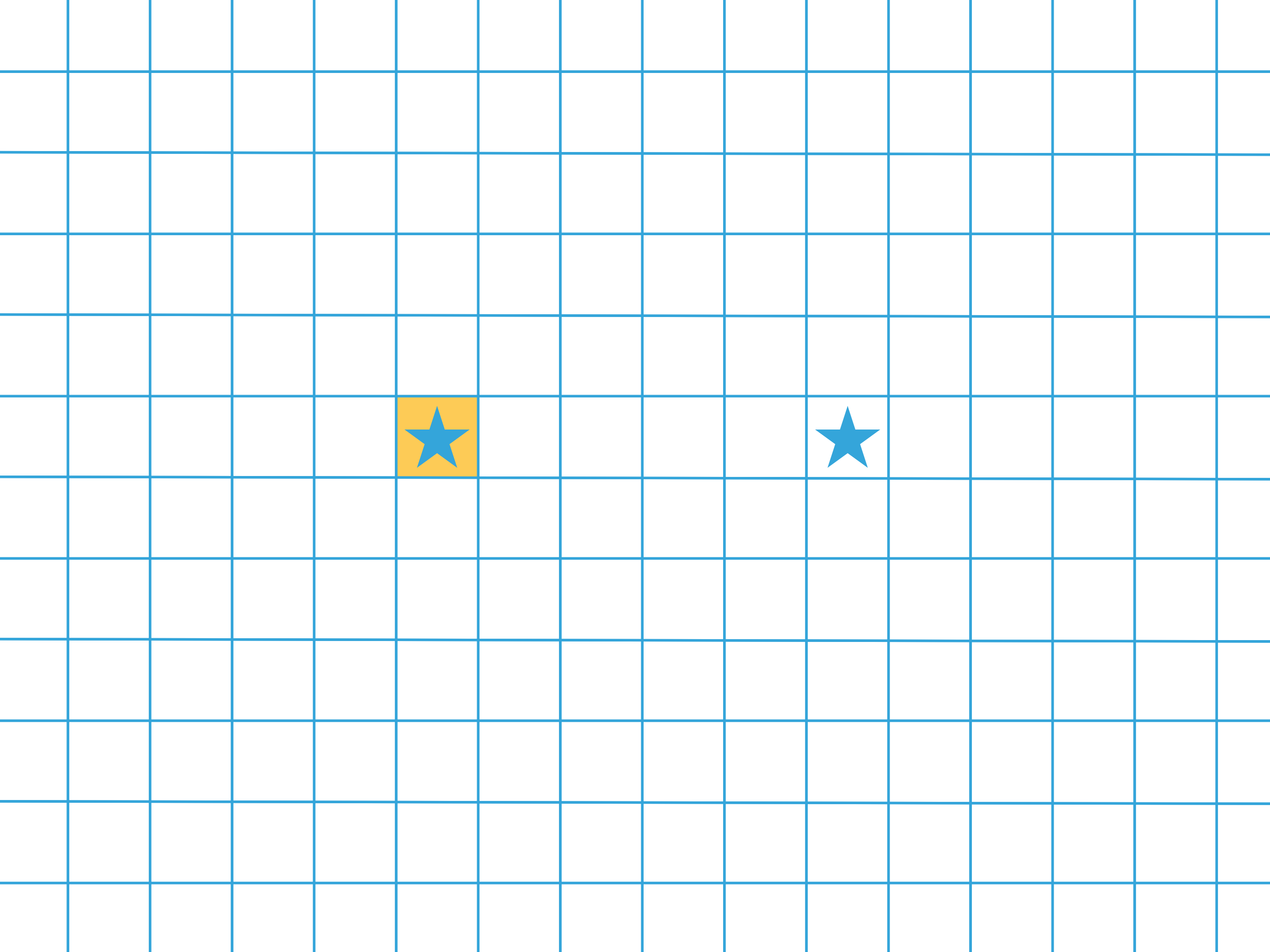


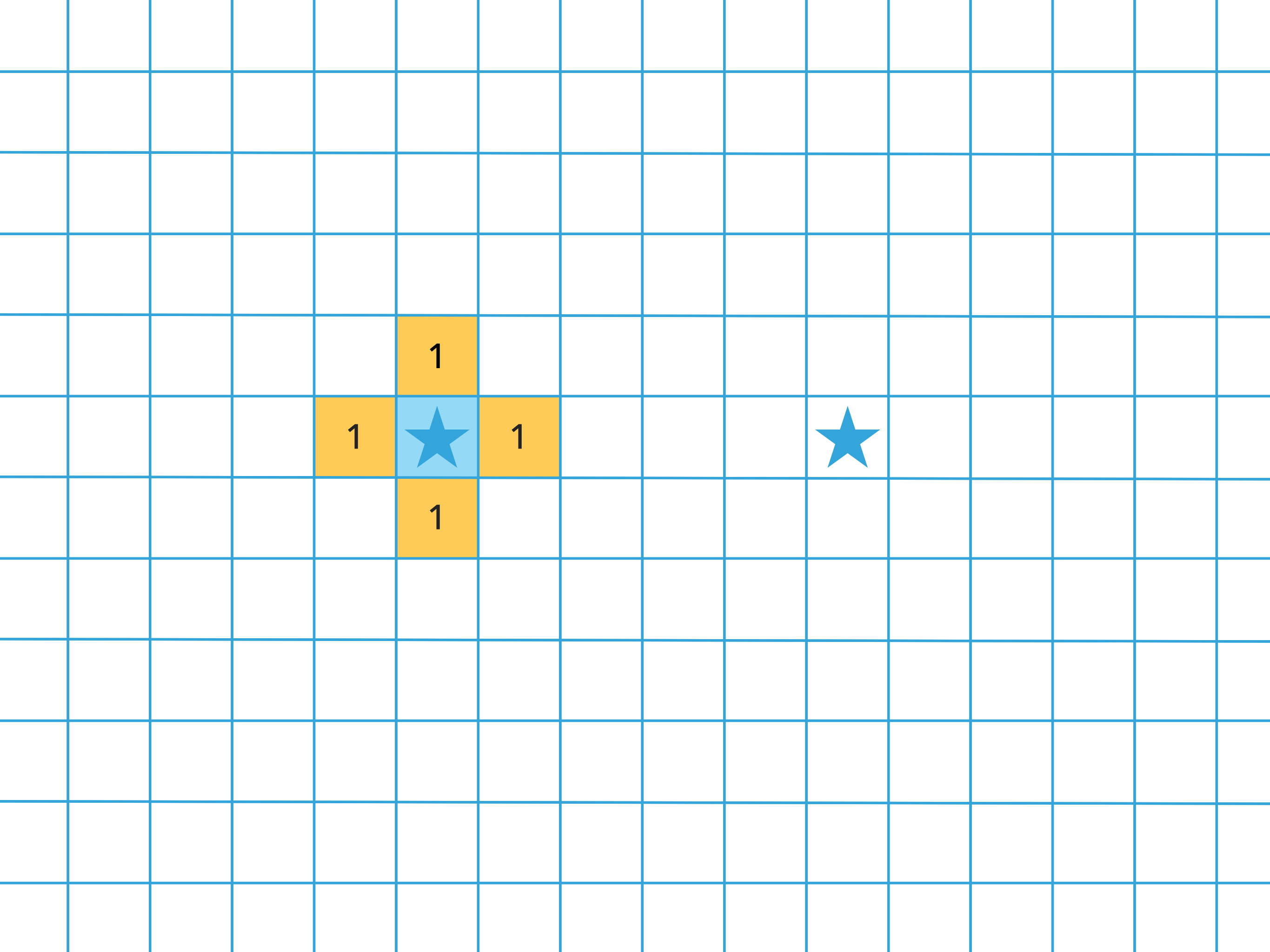
NEGATIVE CYCLES

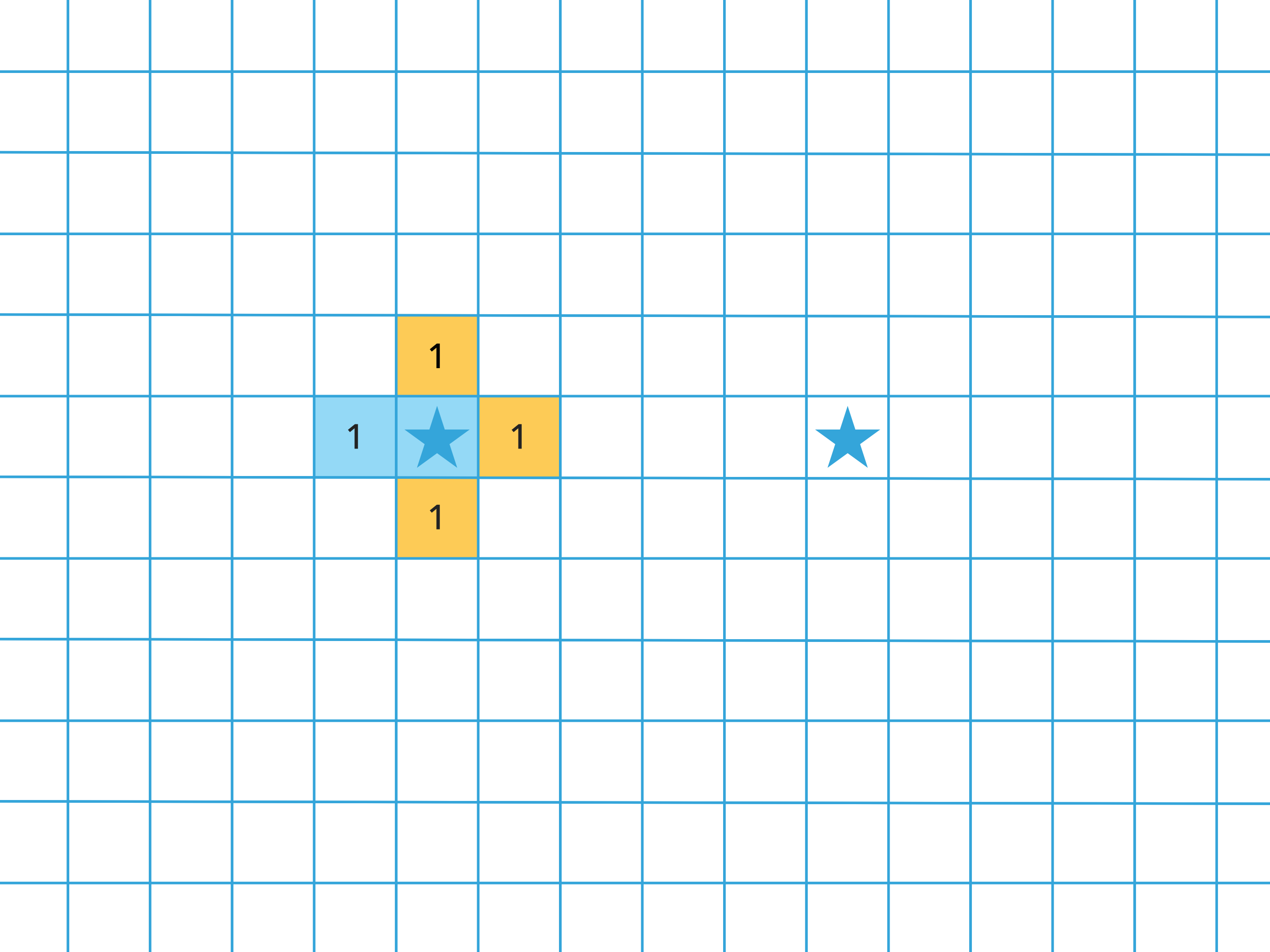


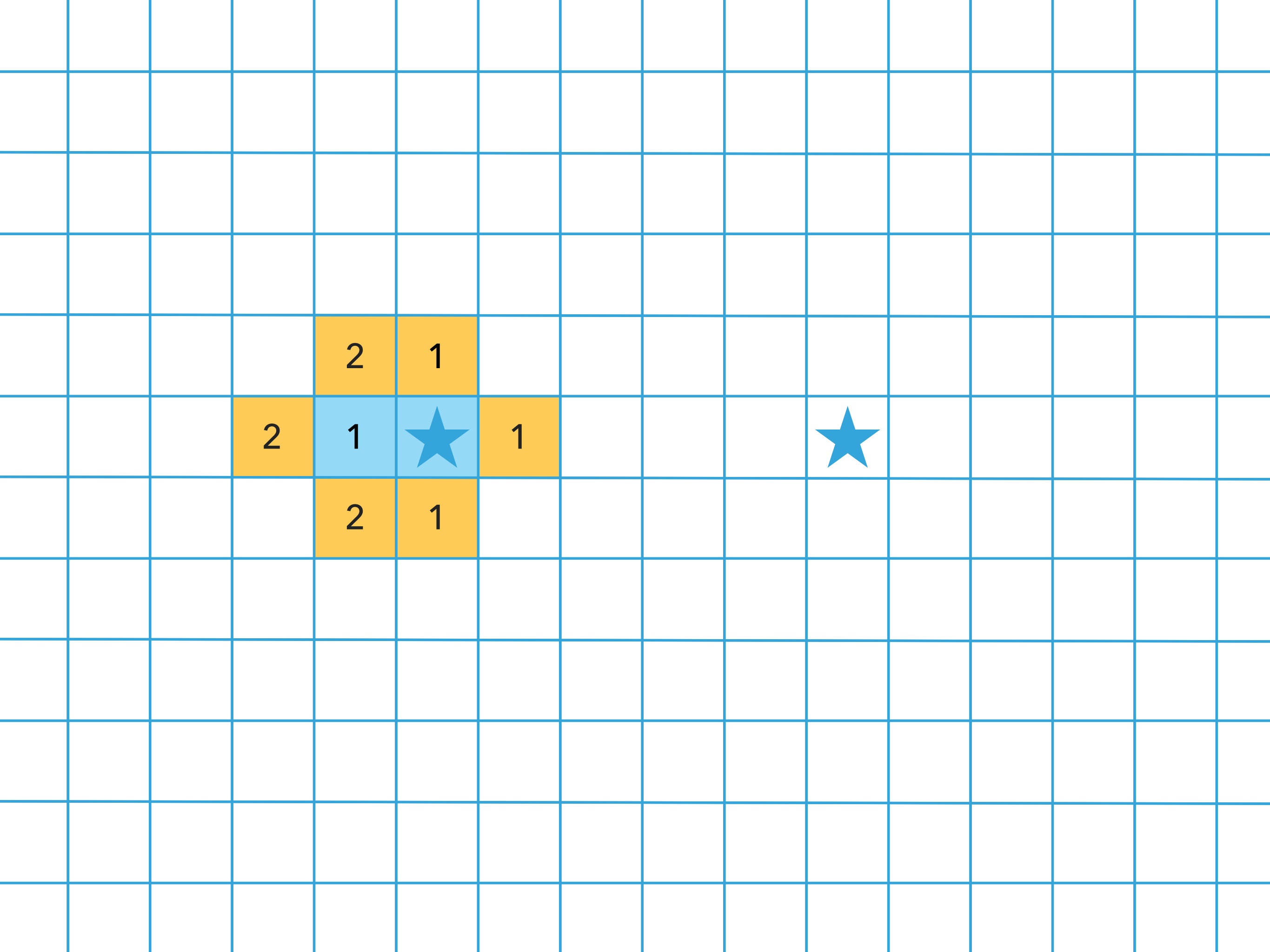












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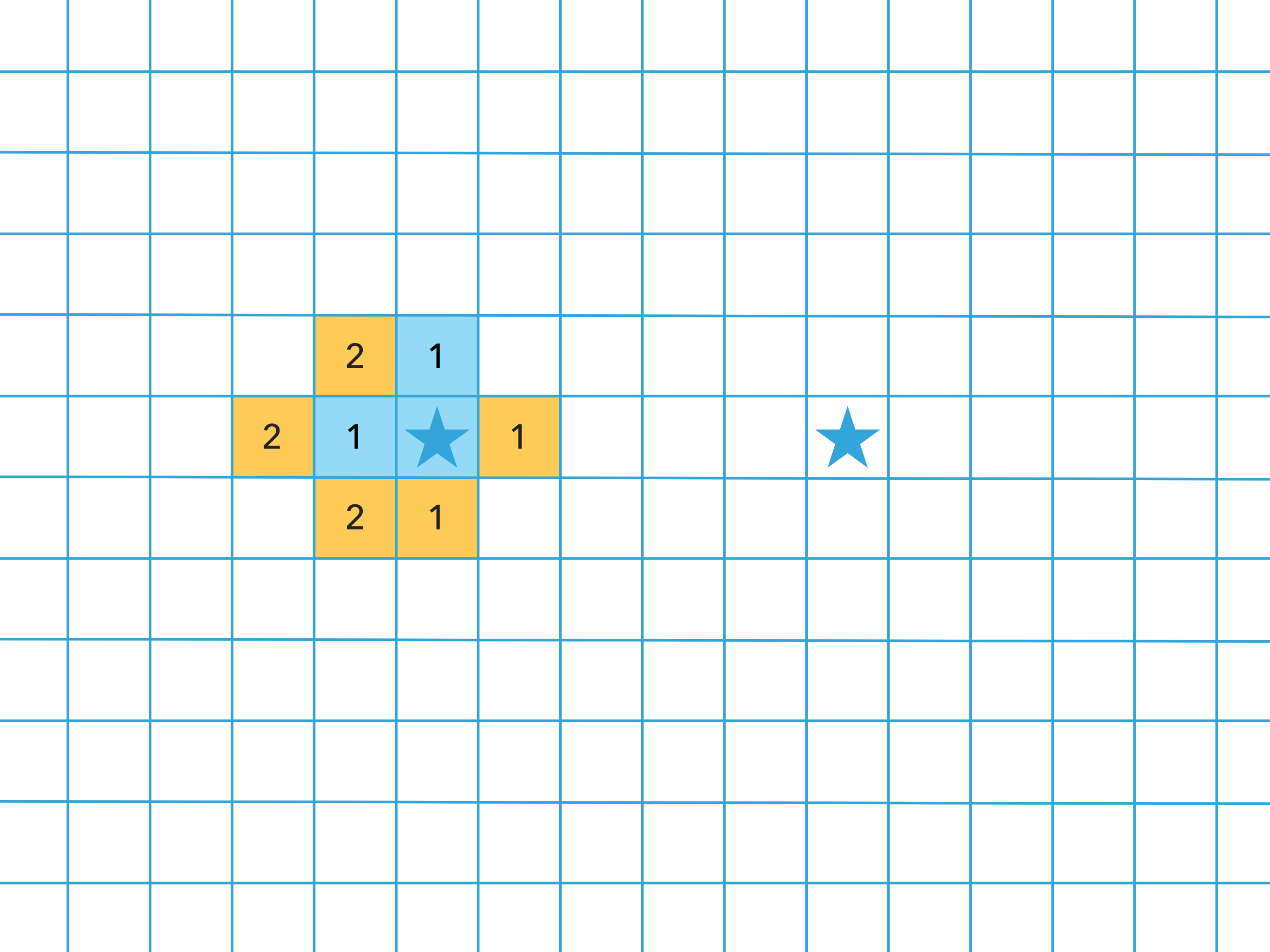
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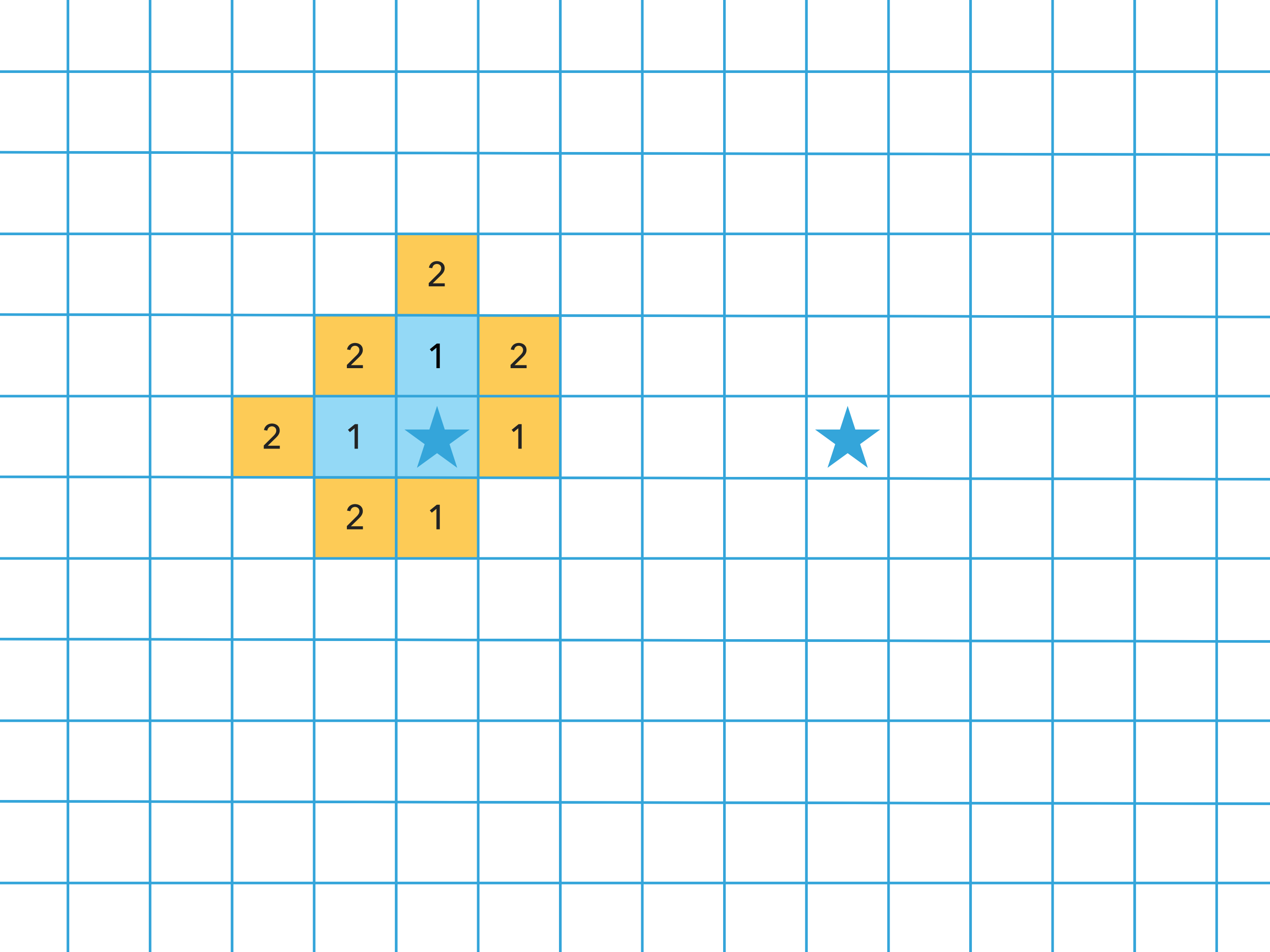
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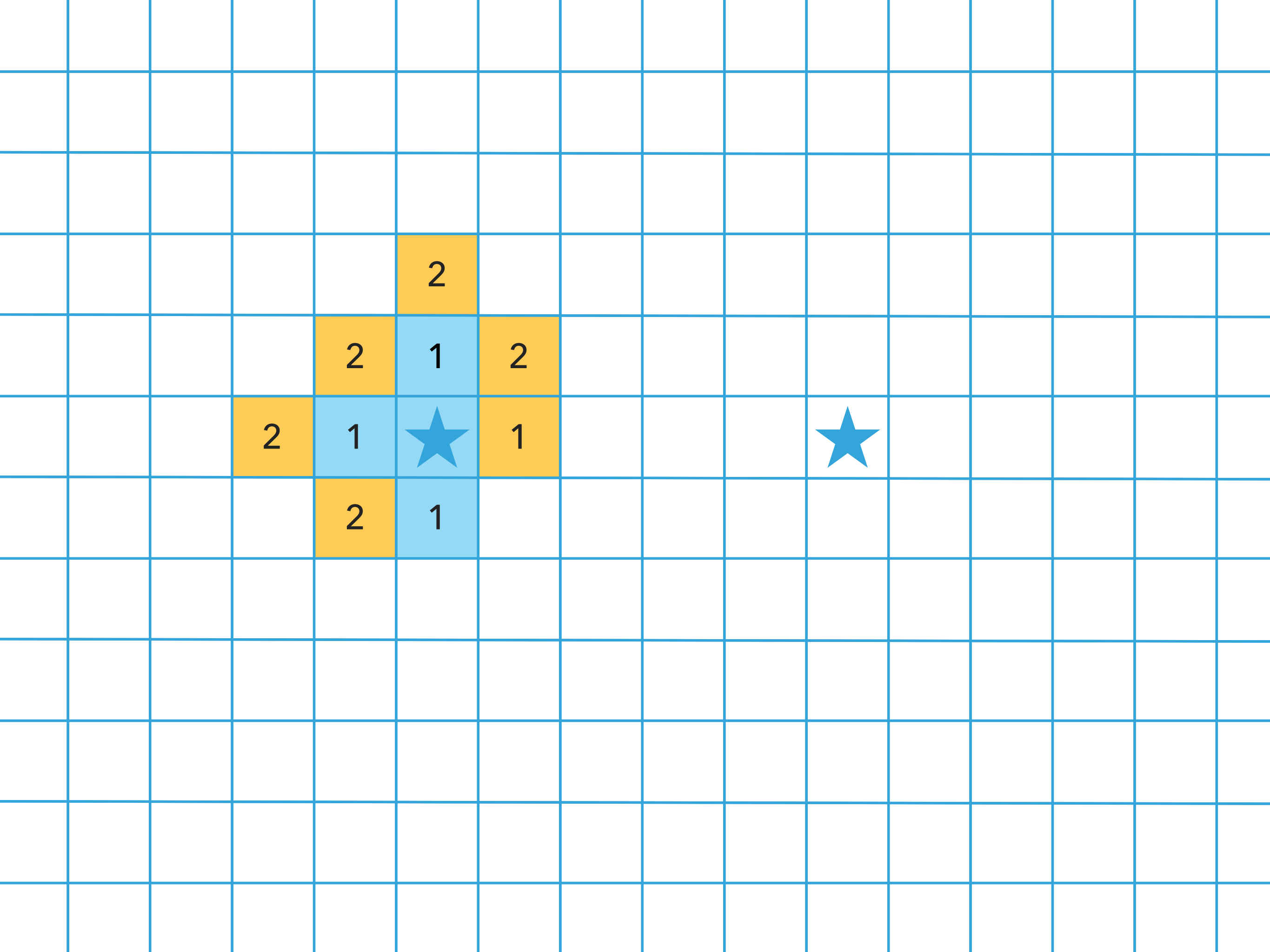
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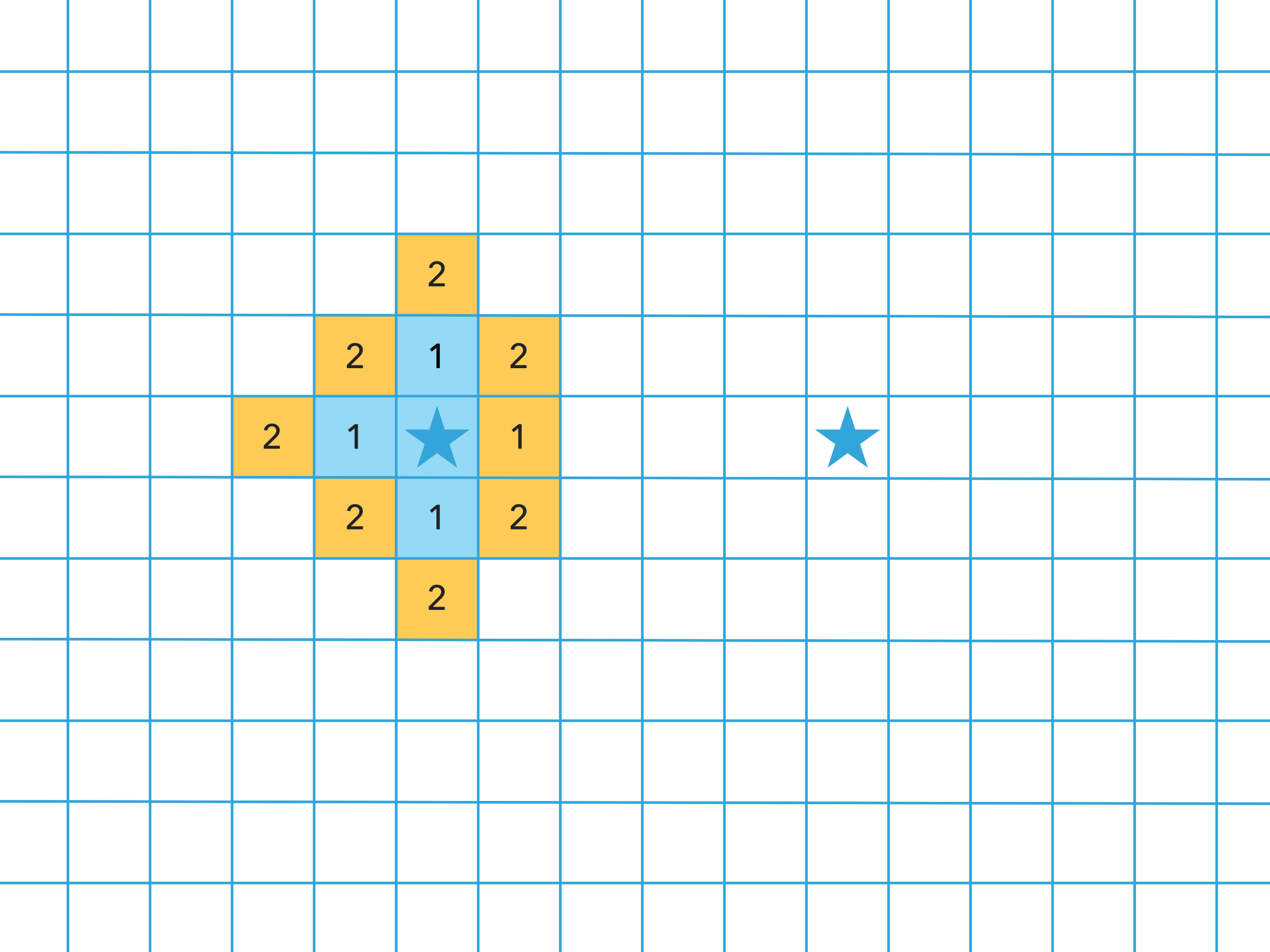
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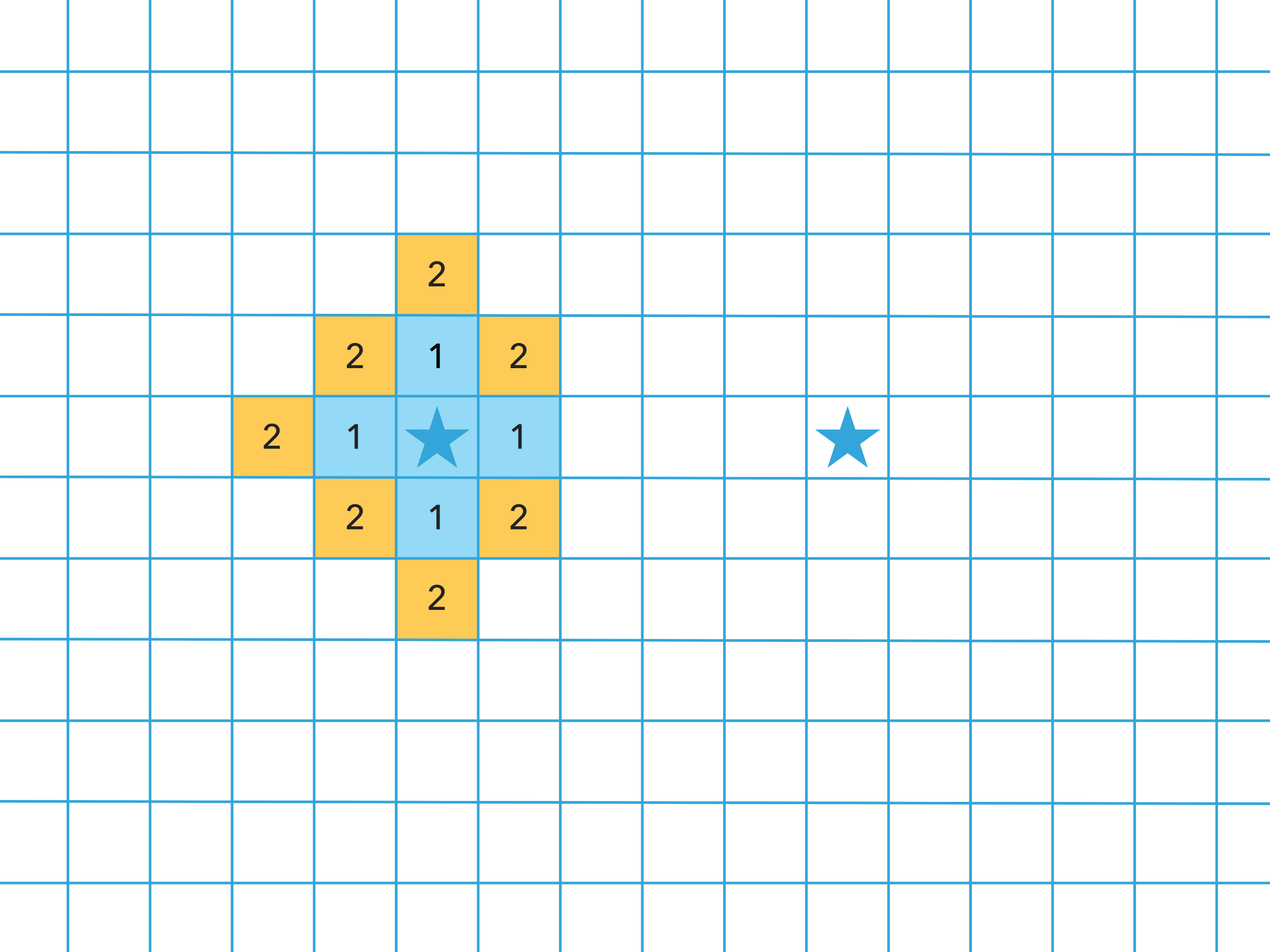


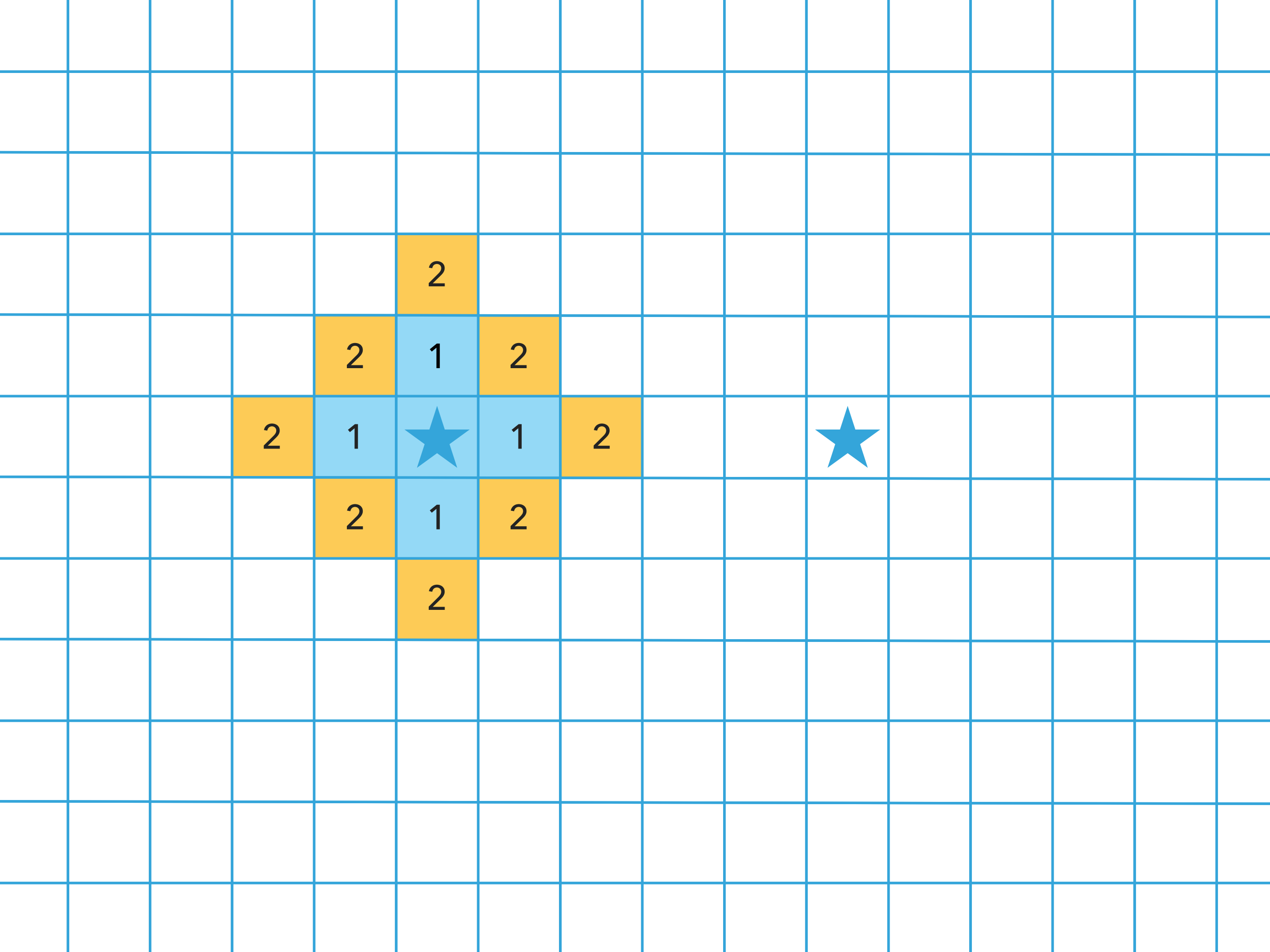
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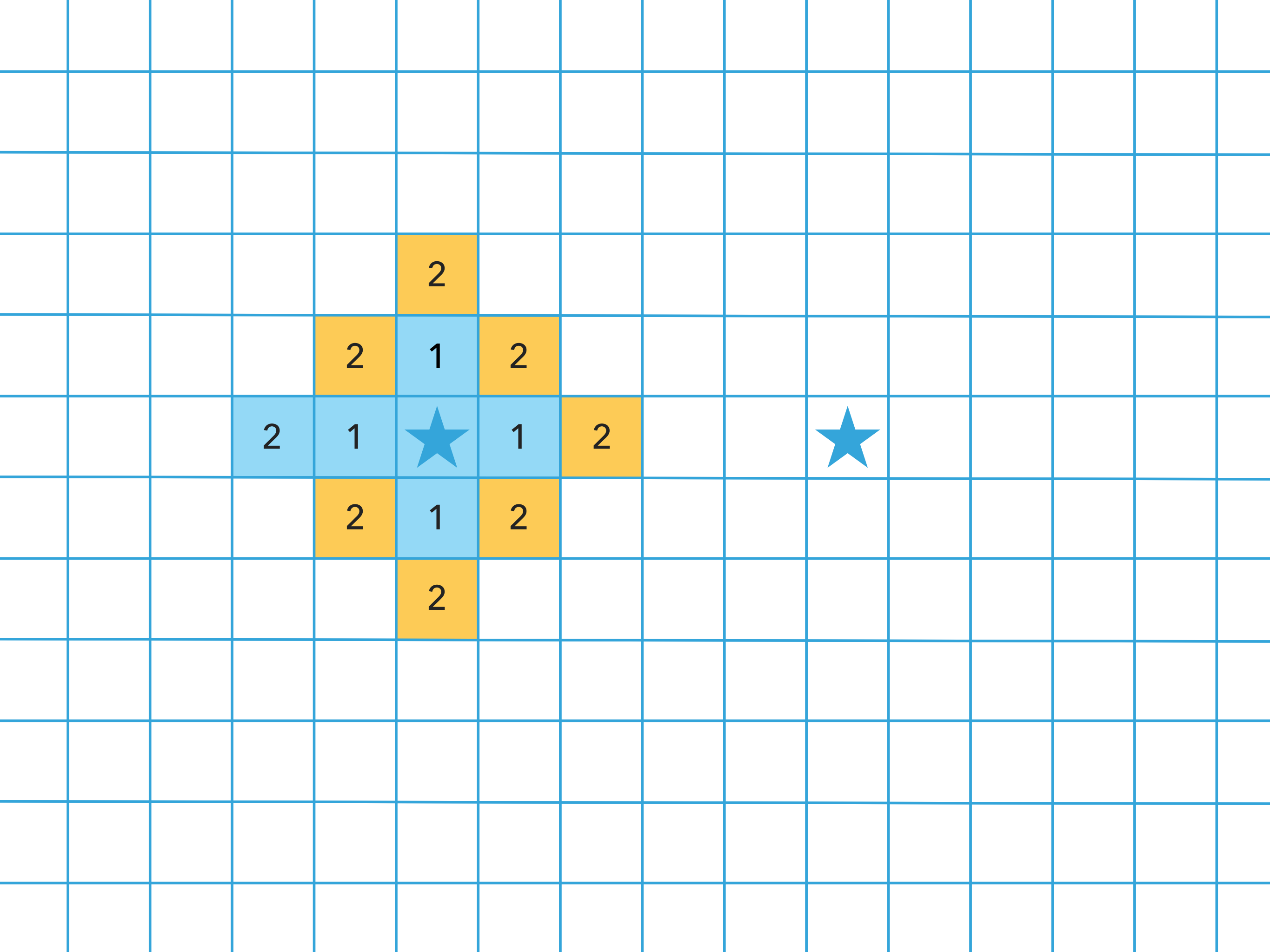


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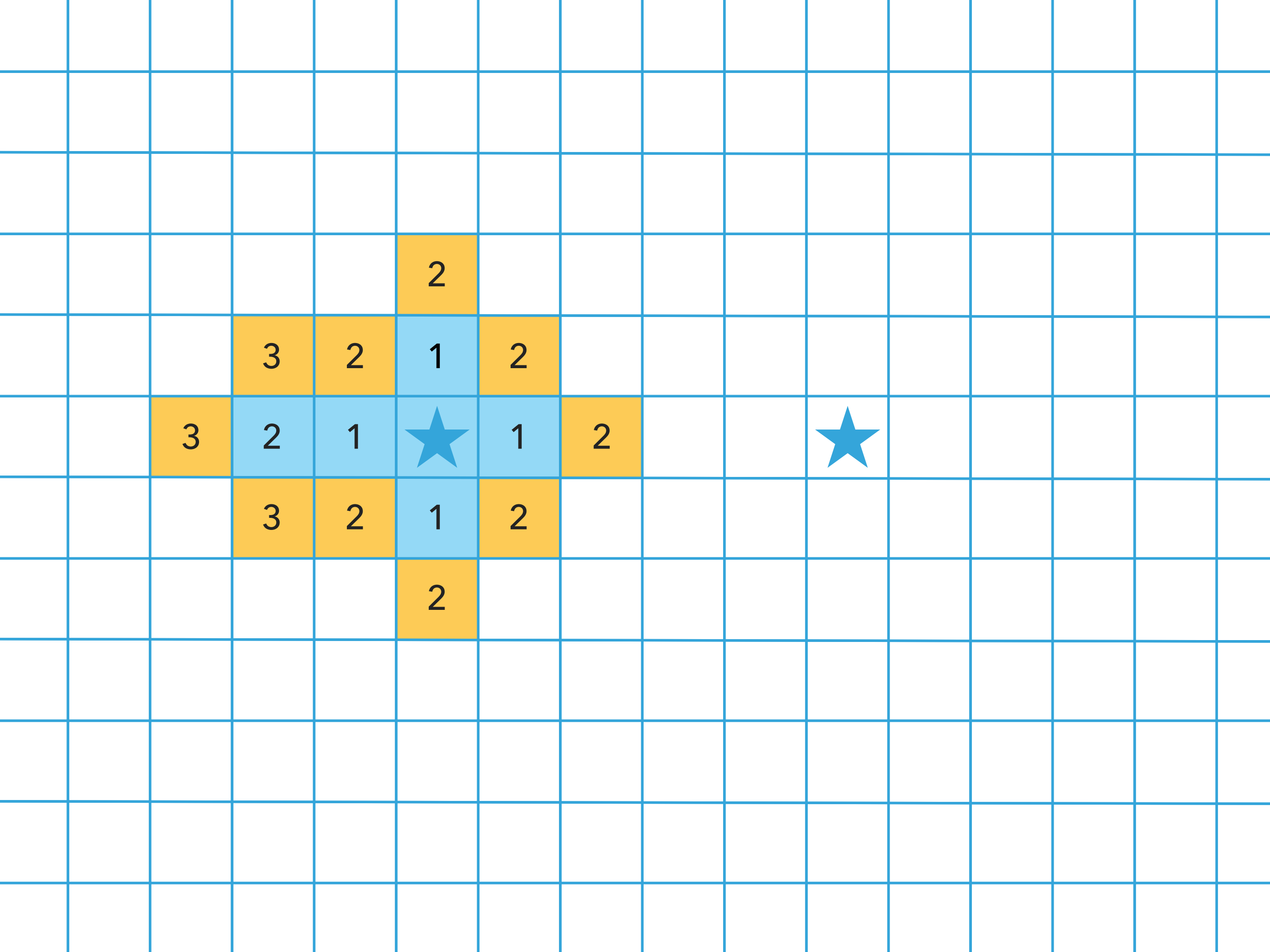


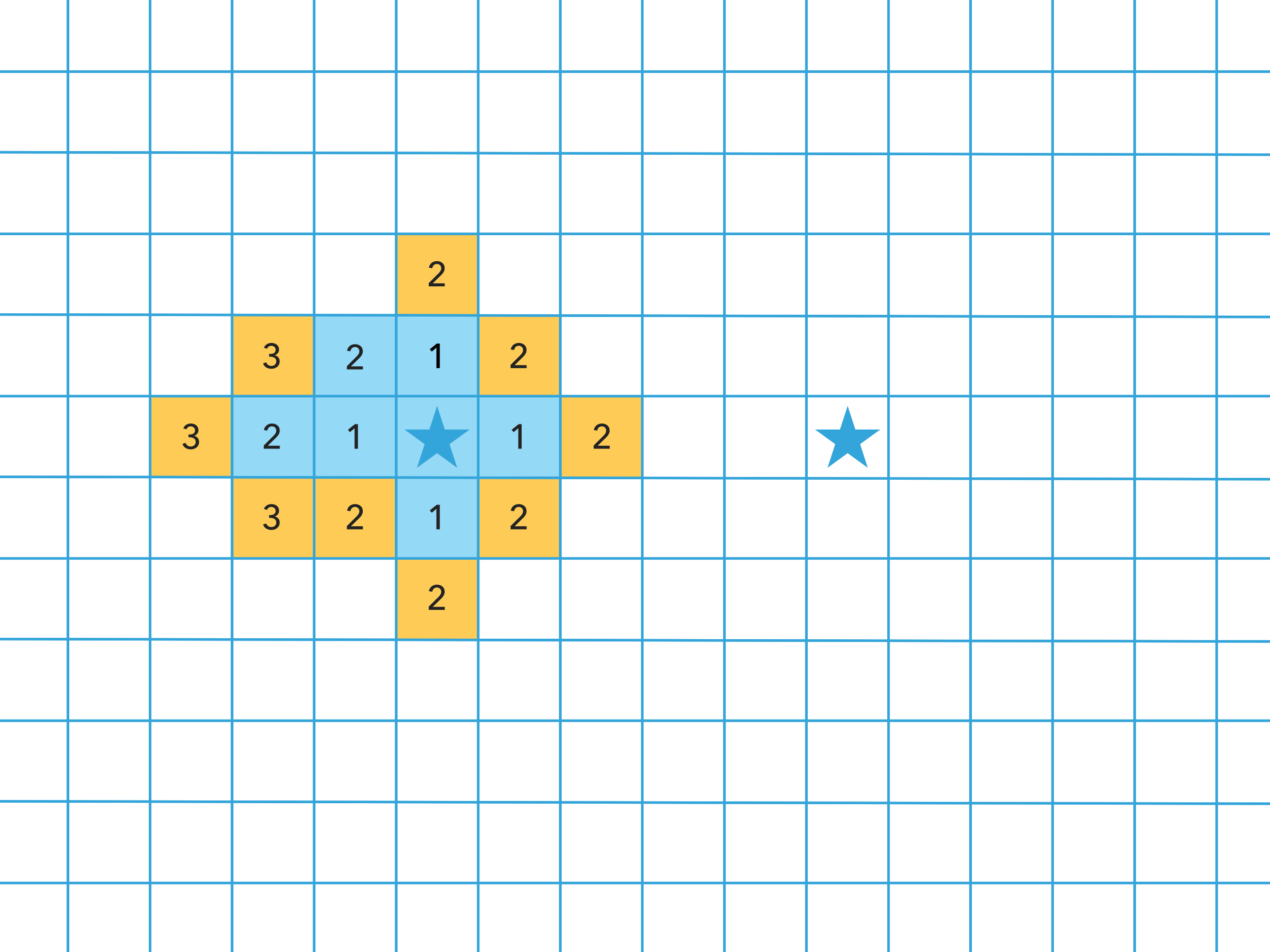
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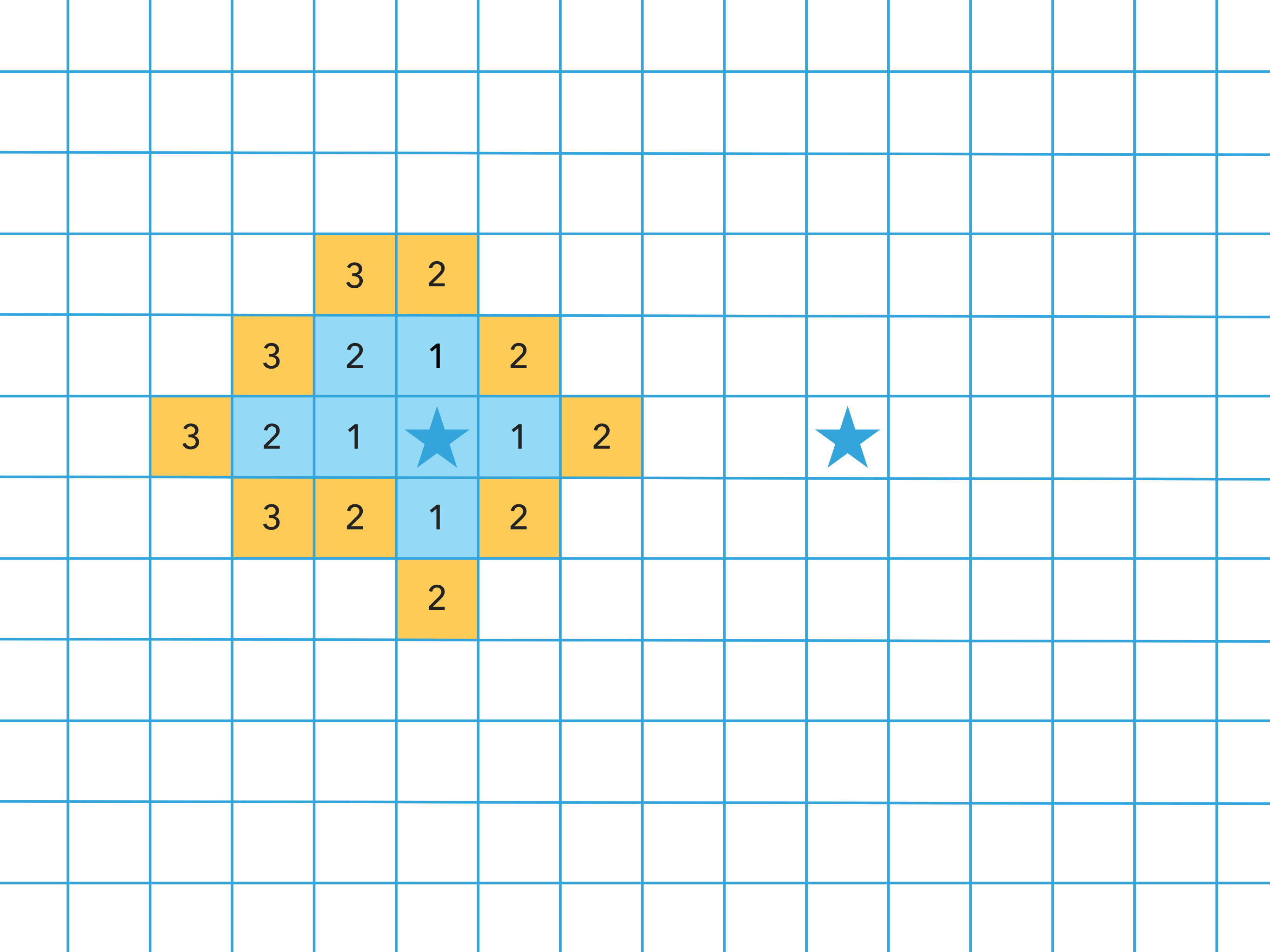
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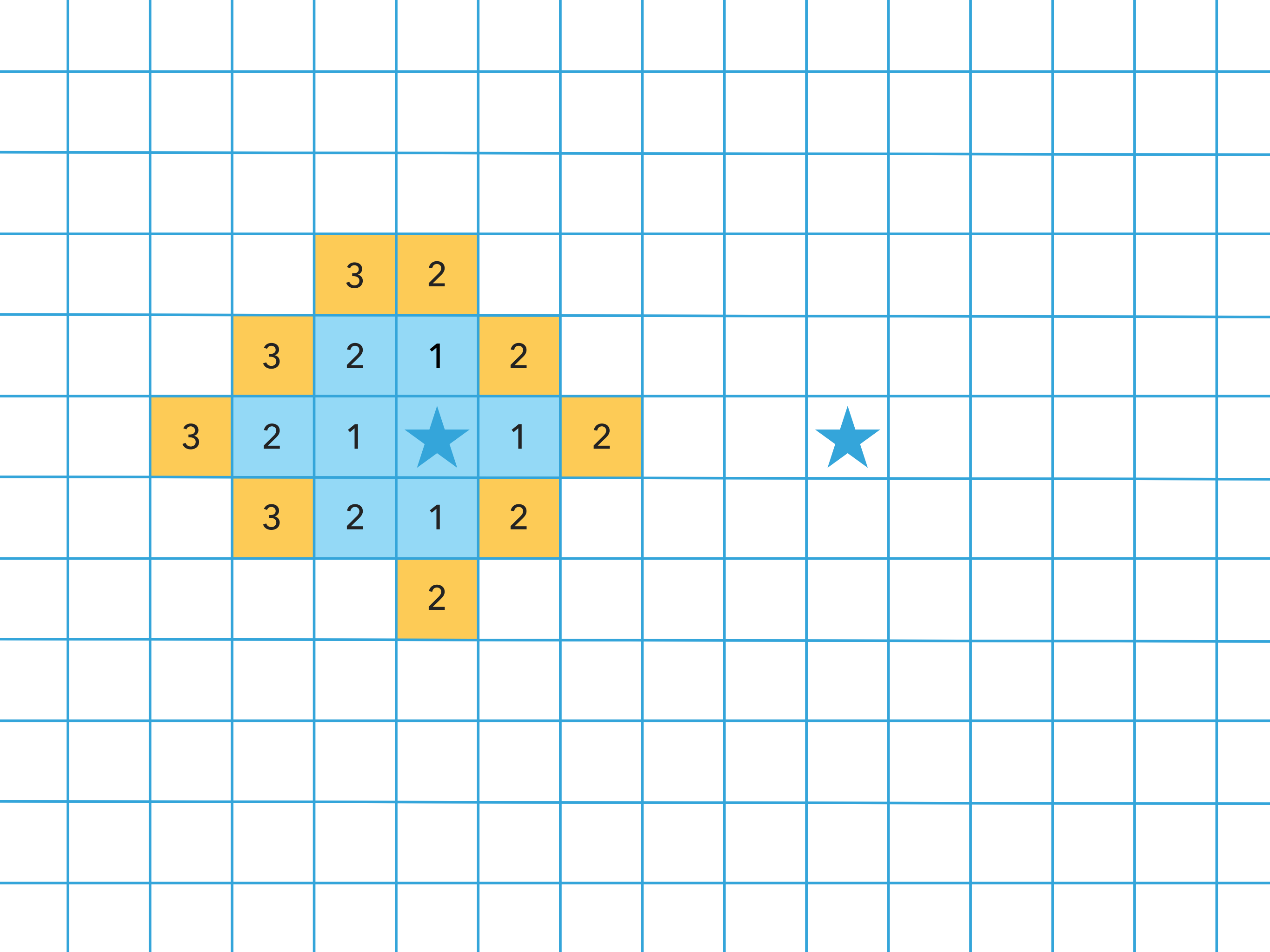
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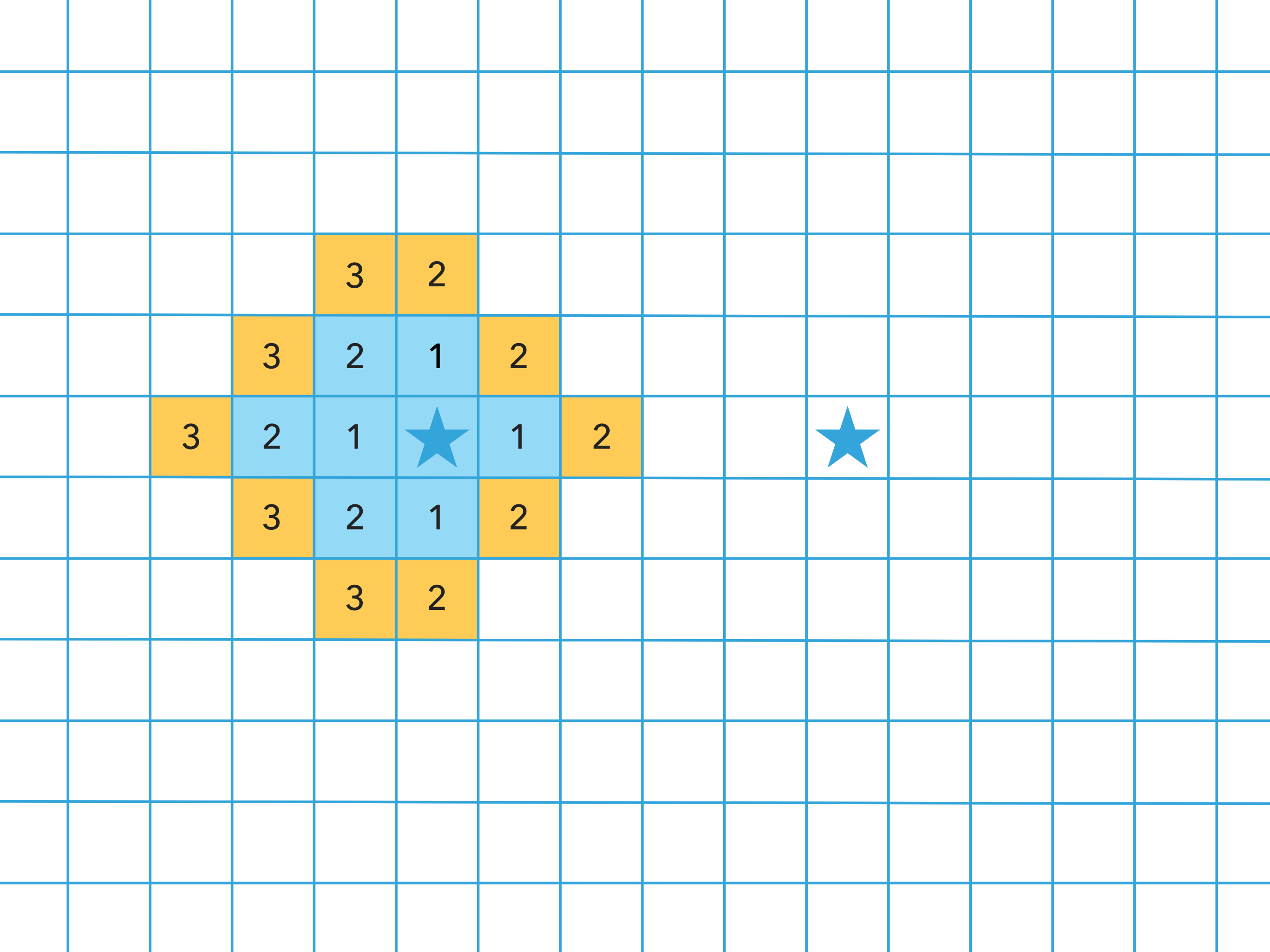
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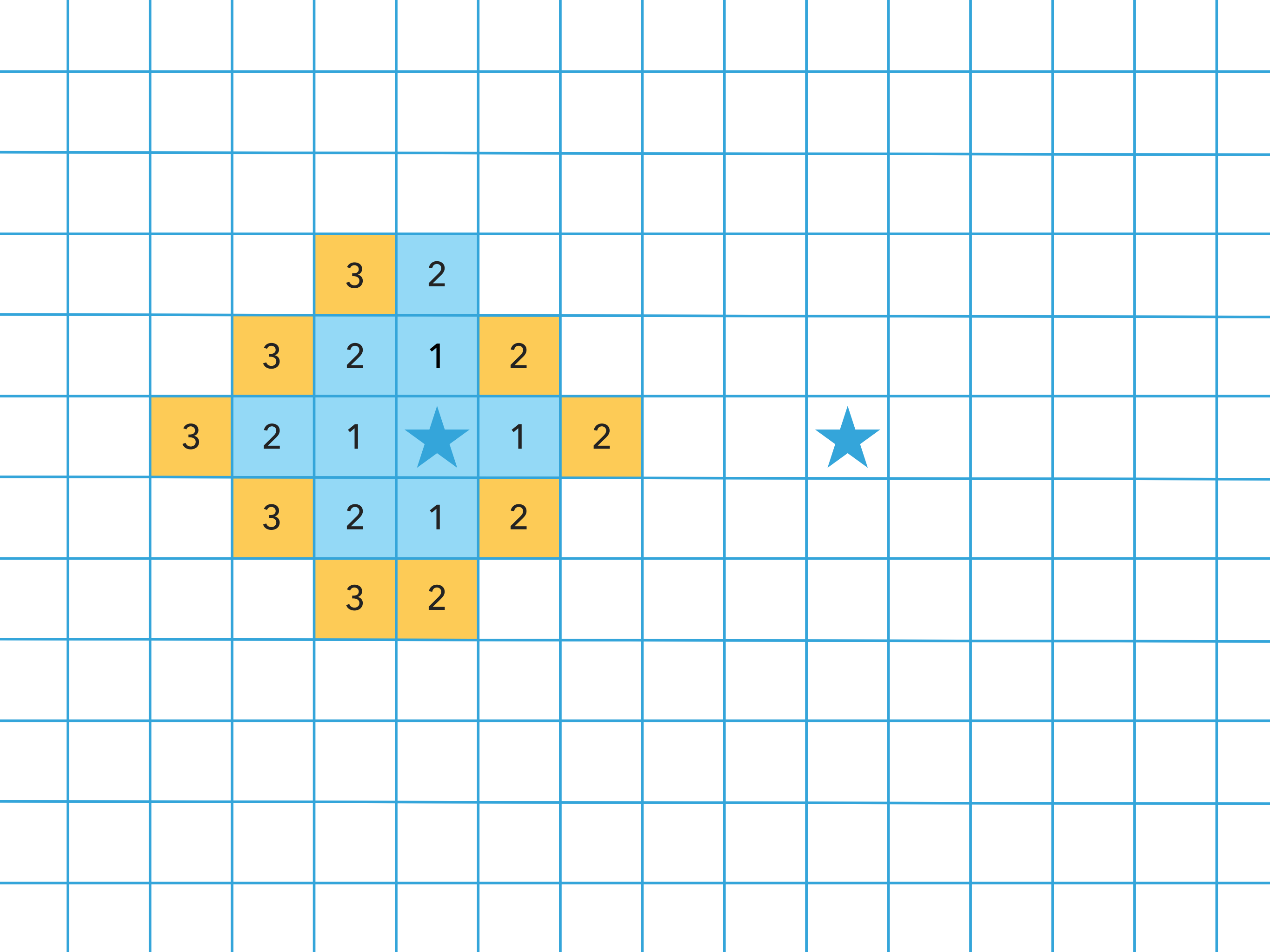


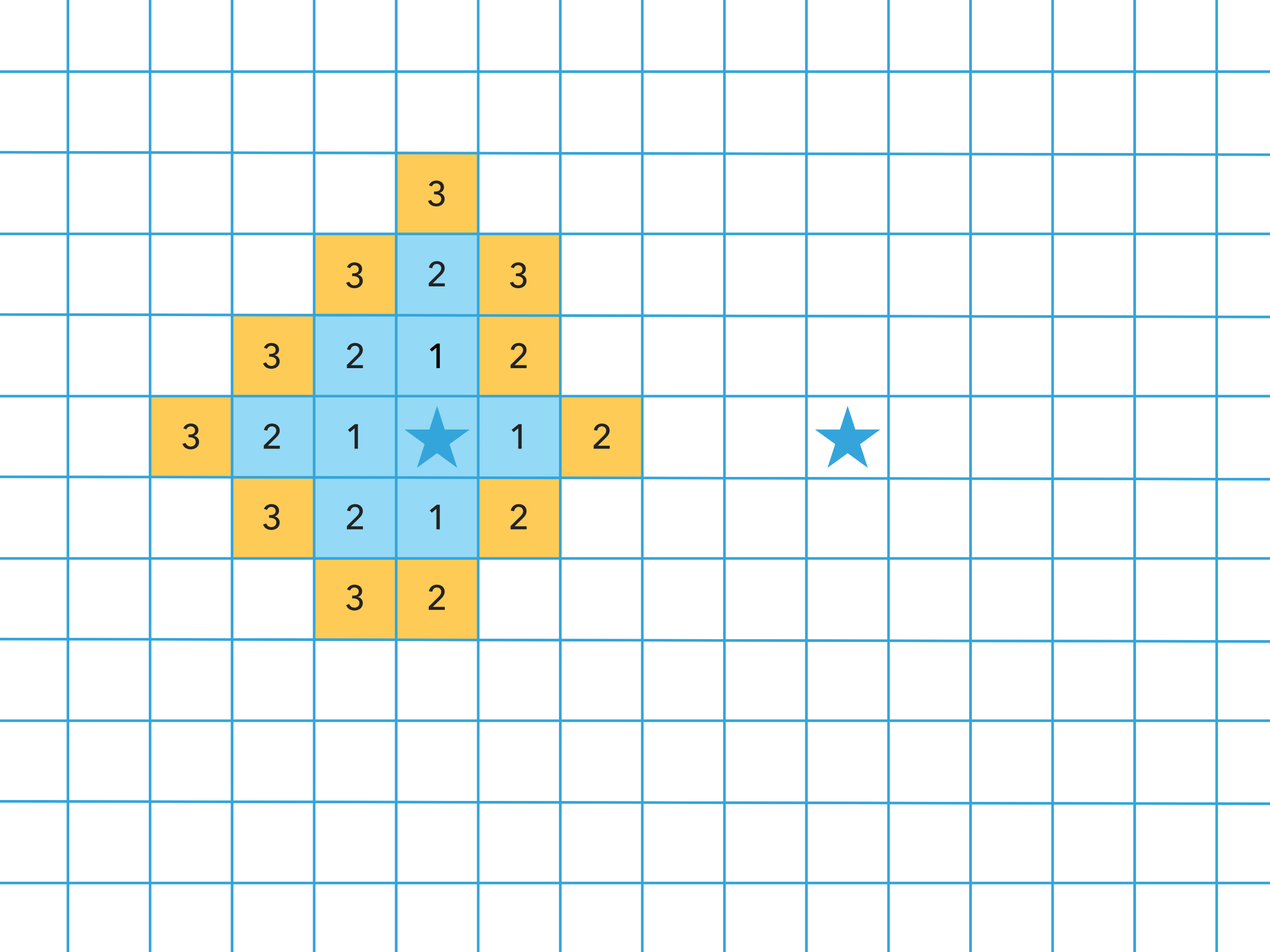


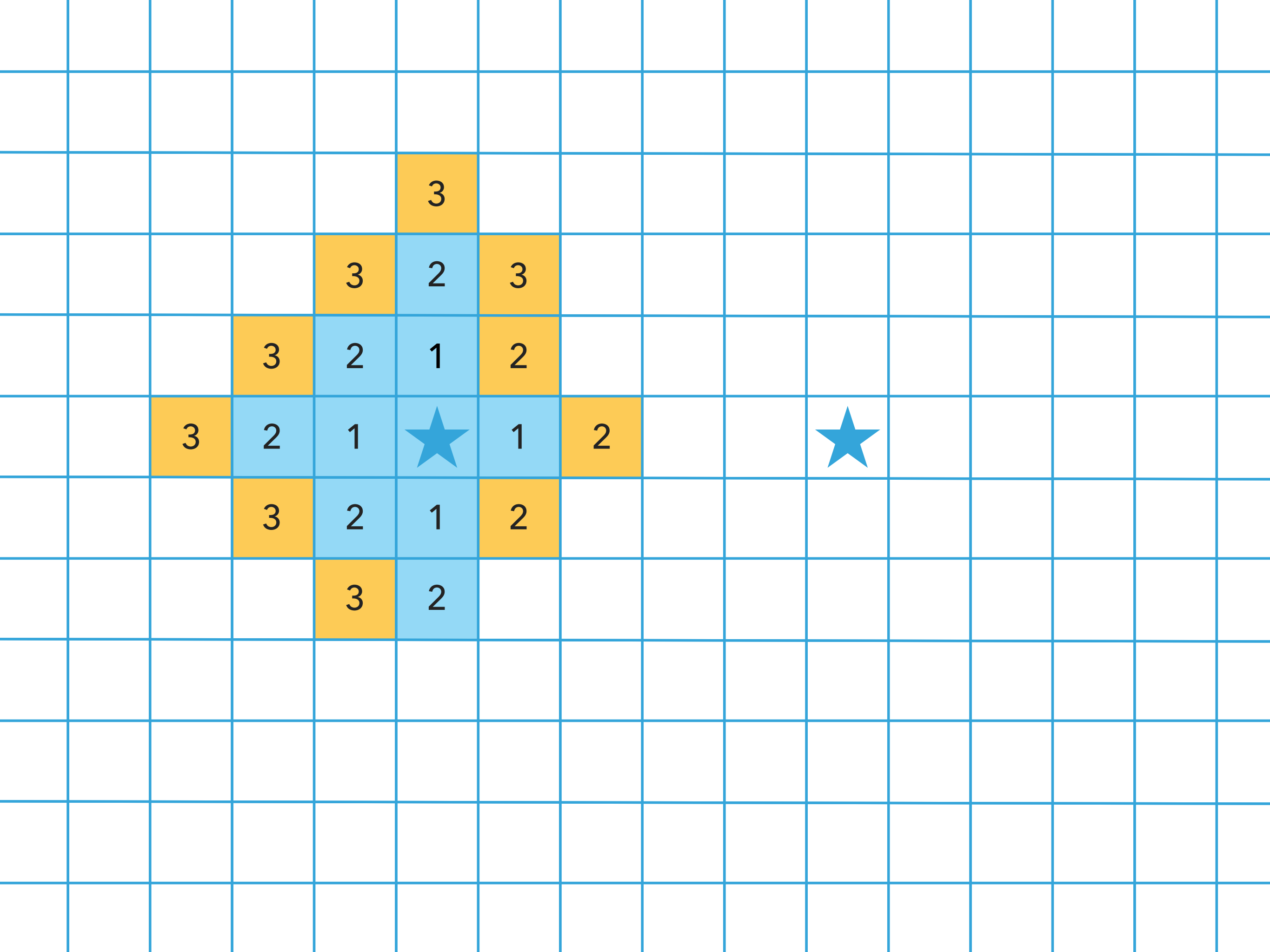


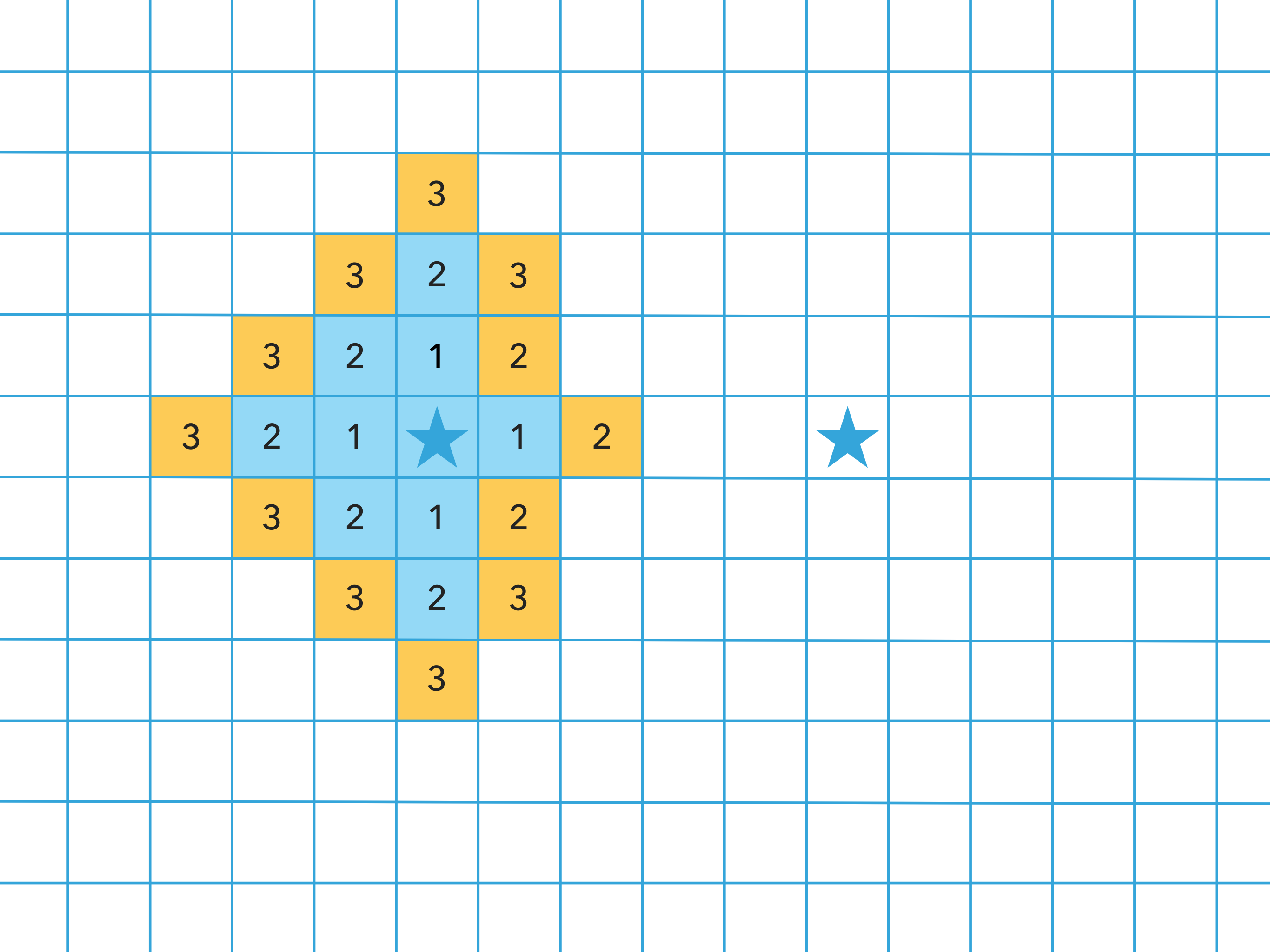


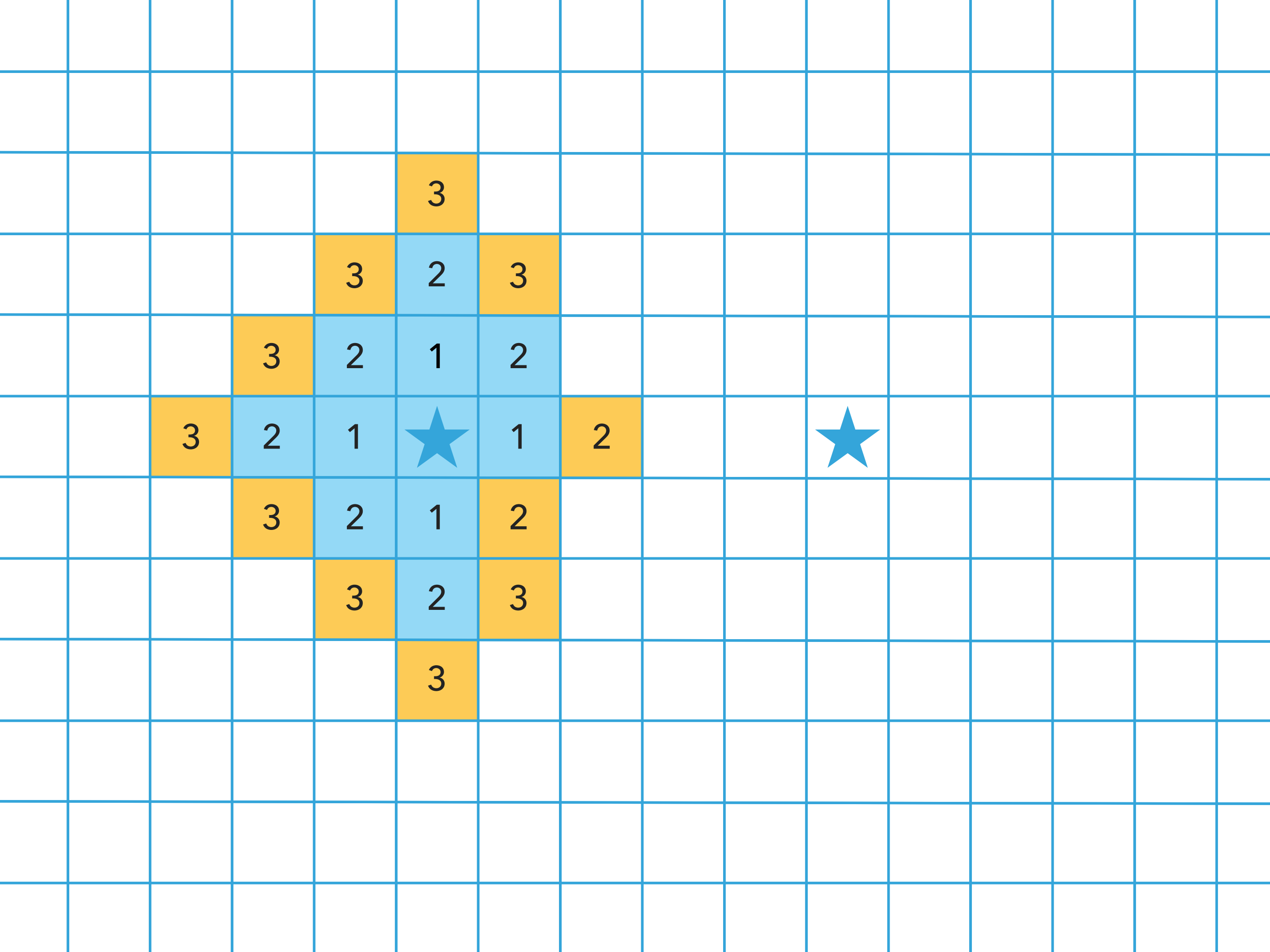


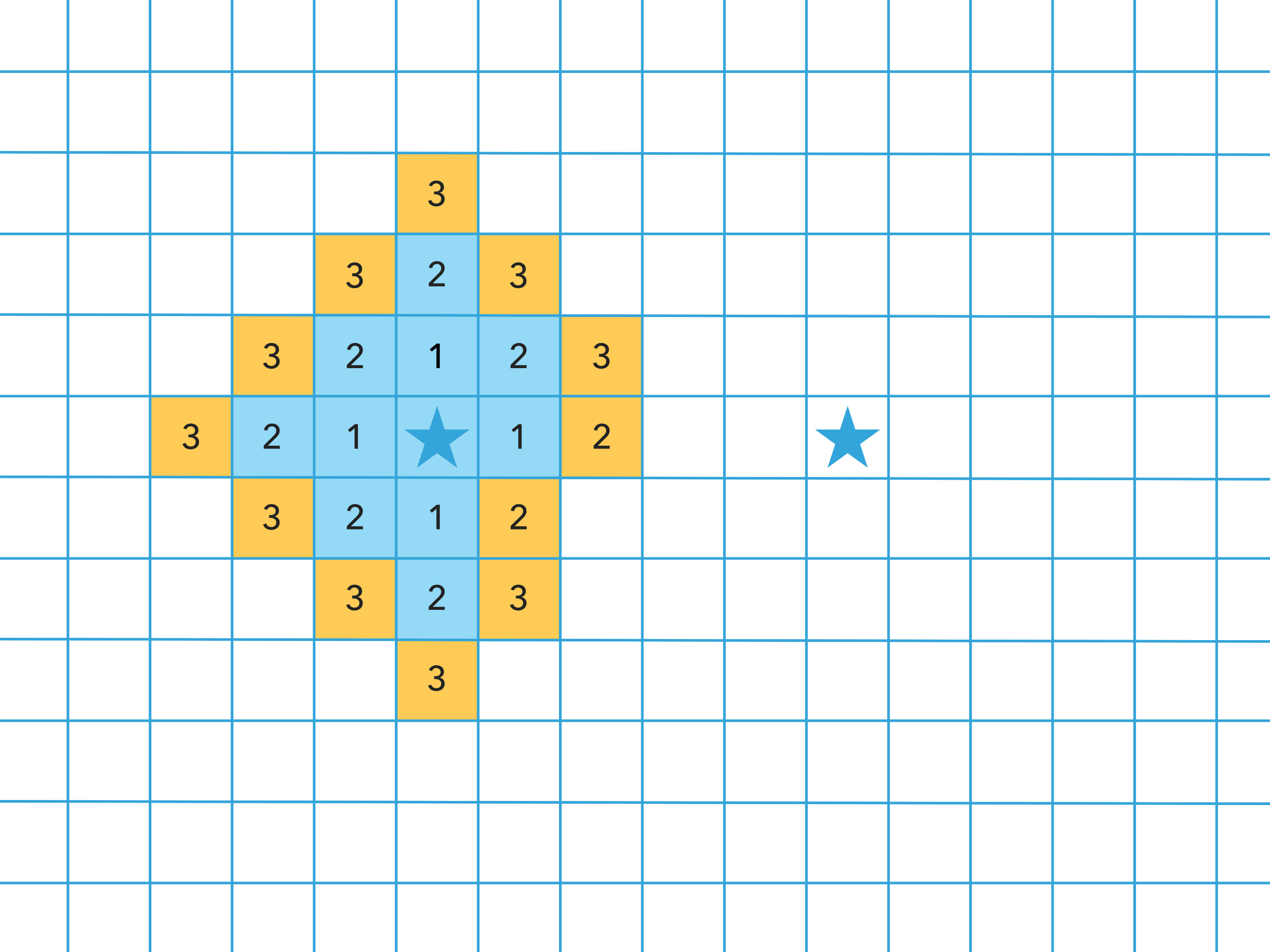


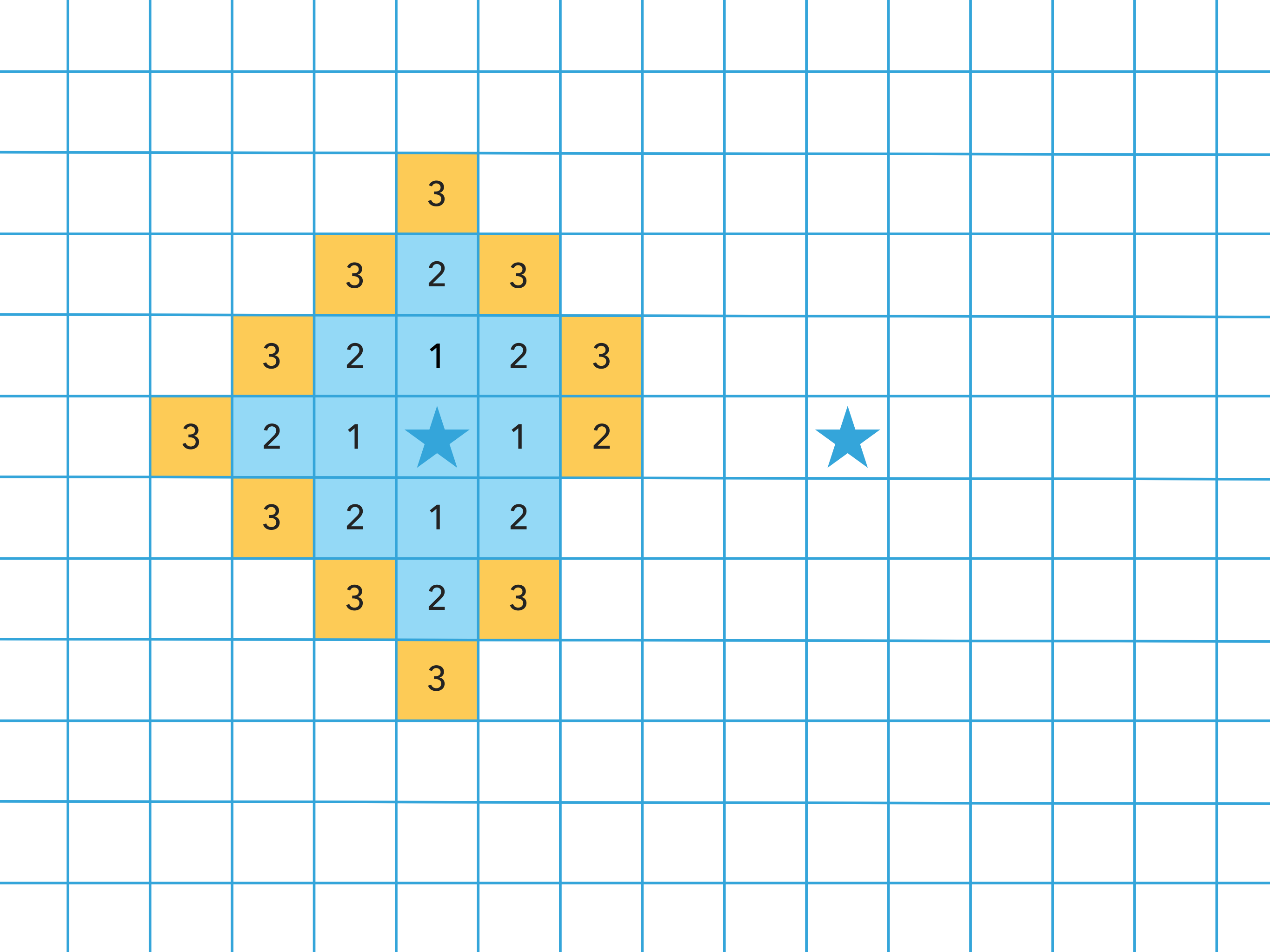


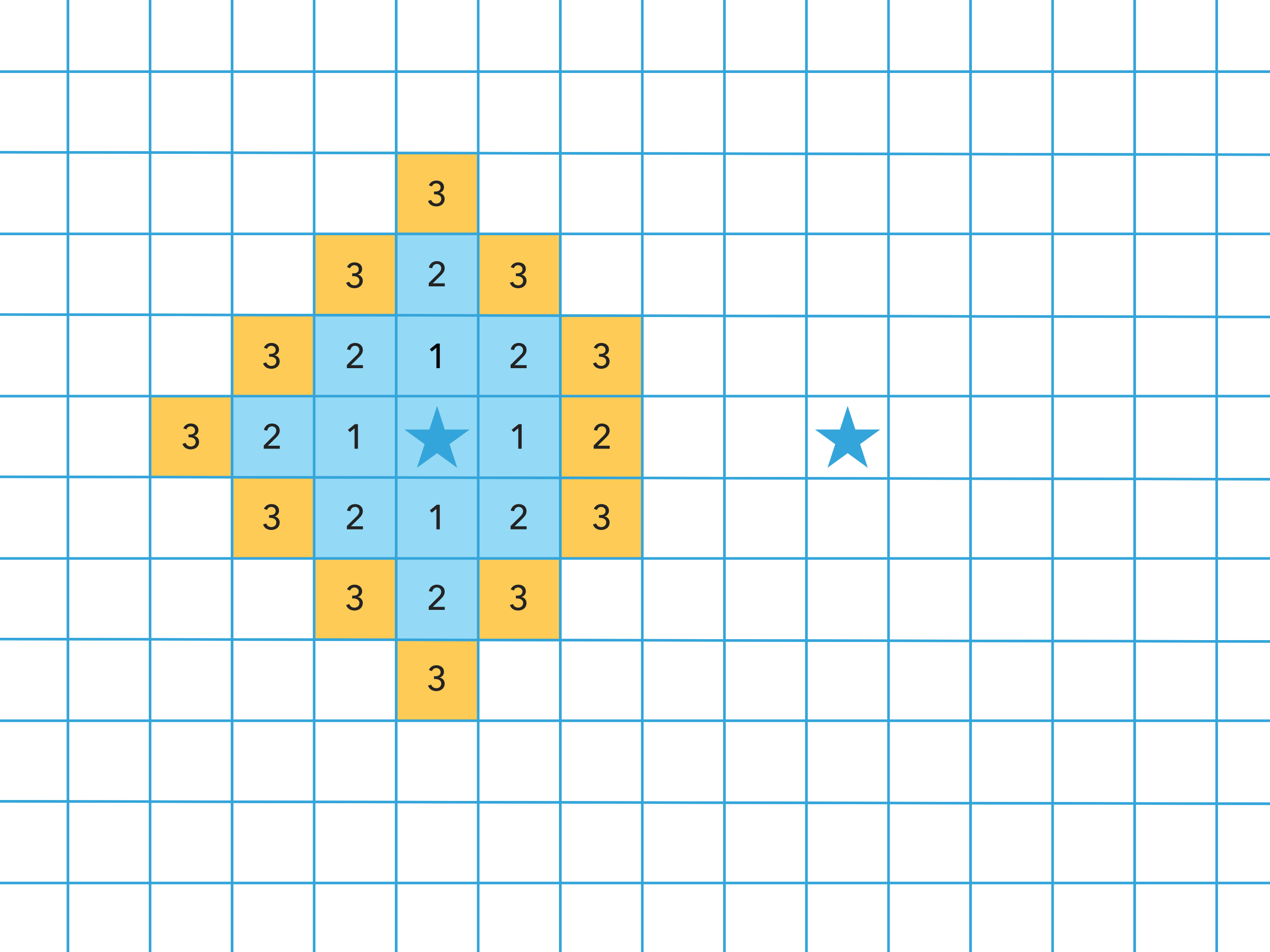


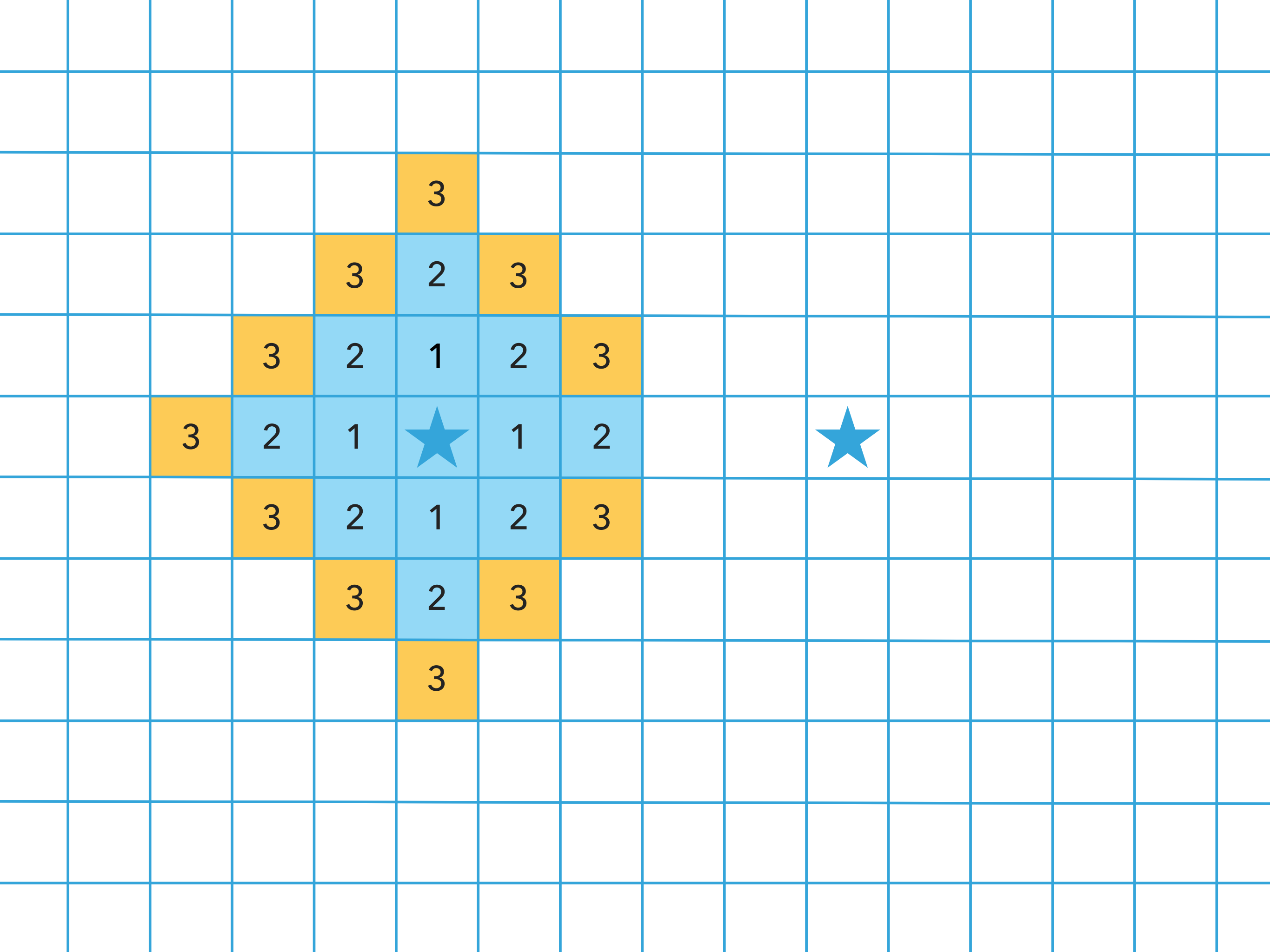


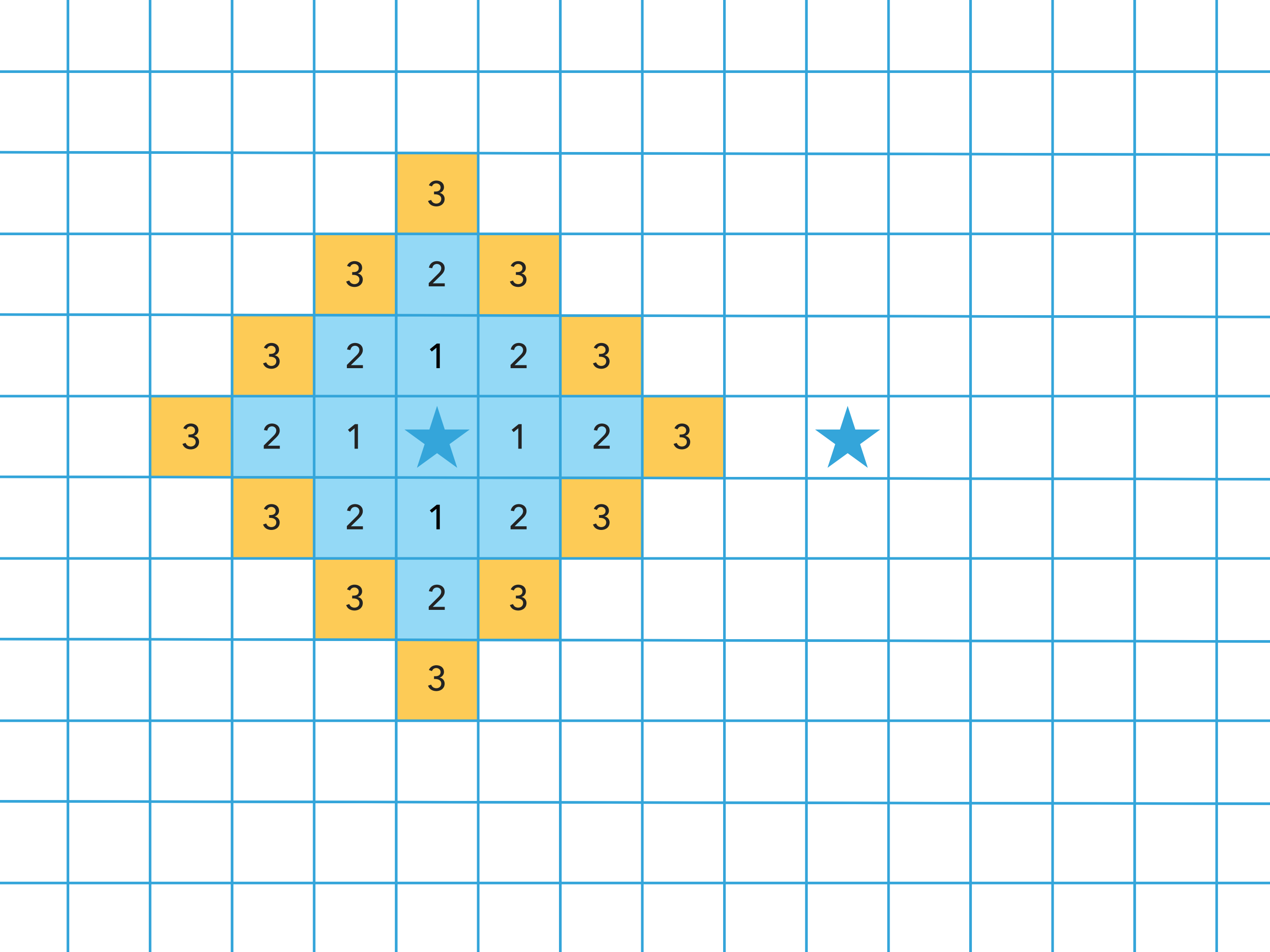


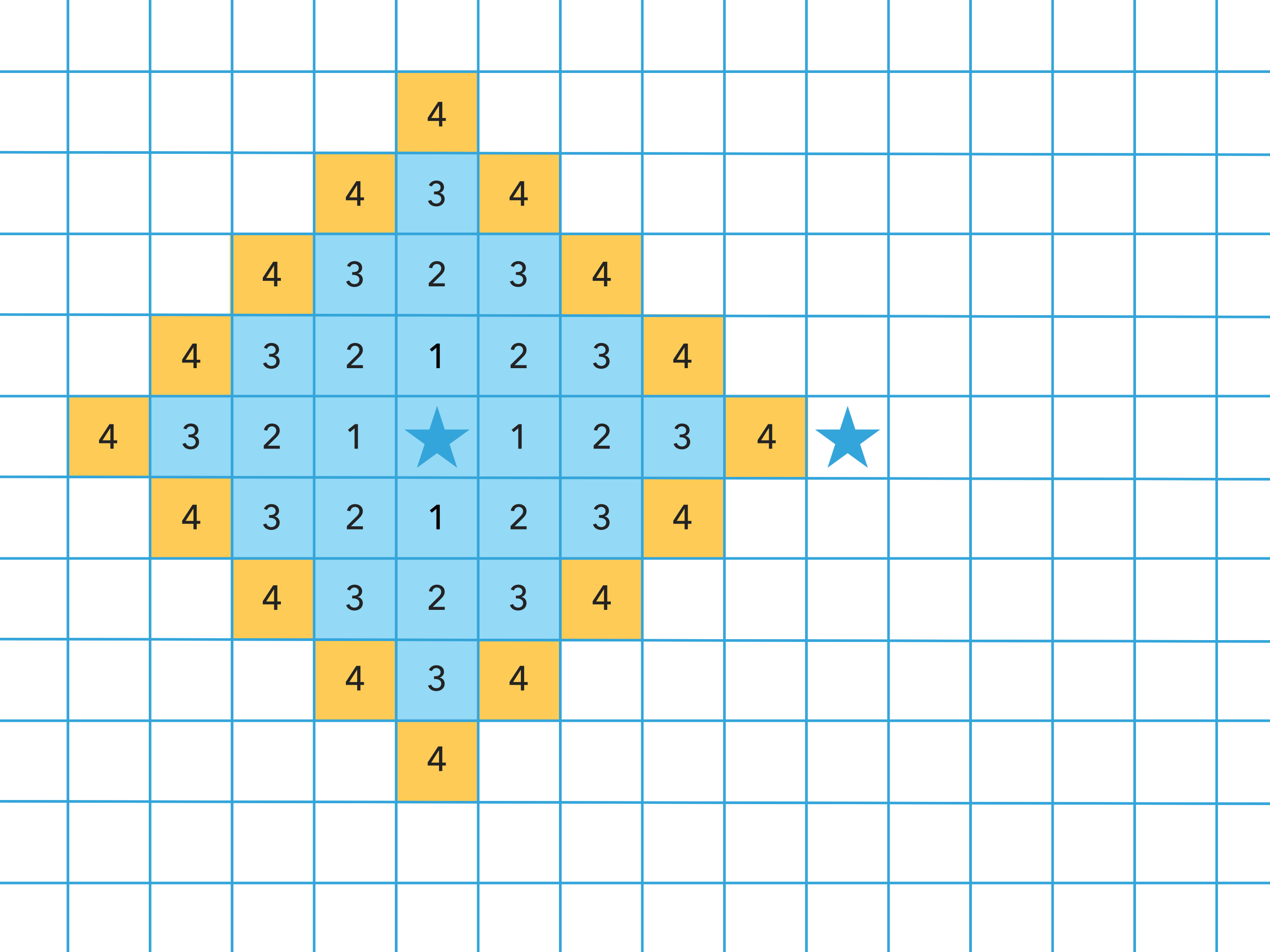


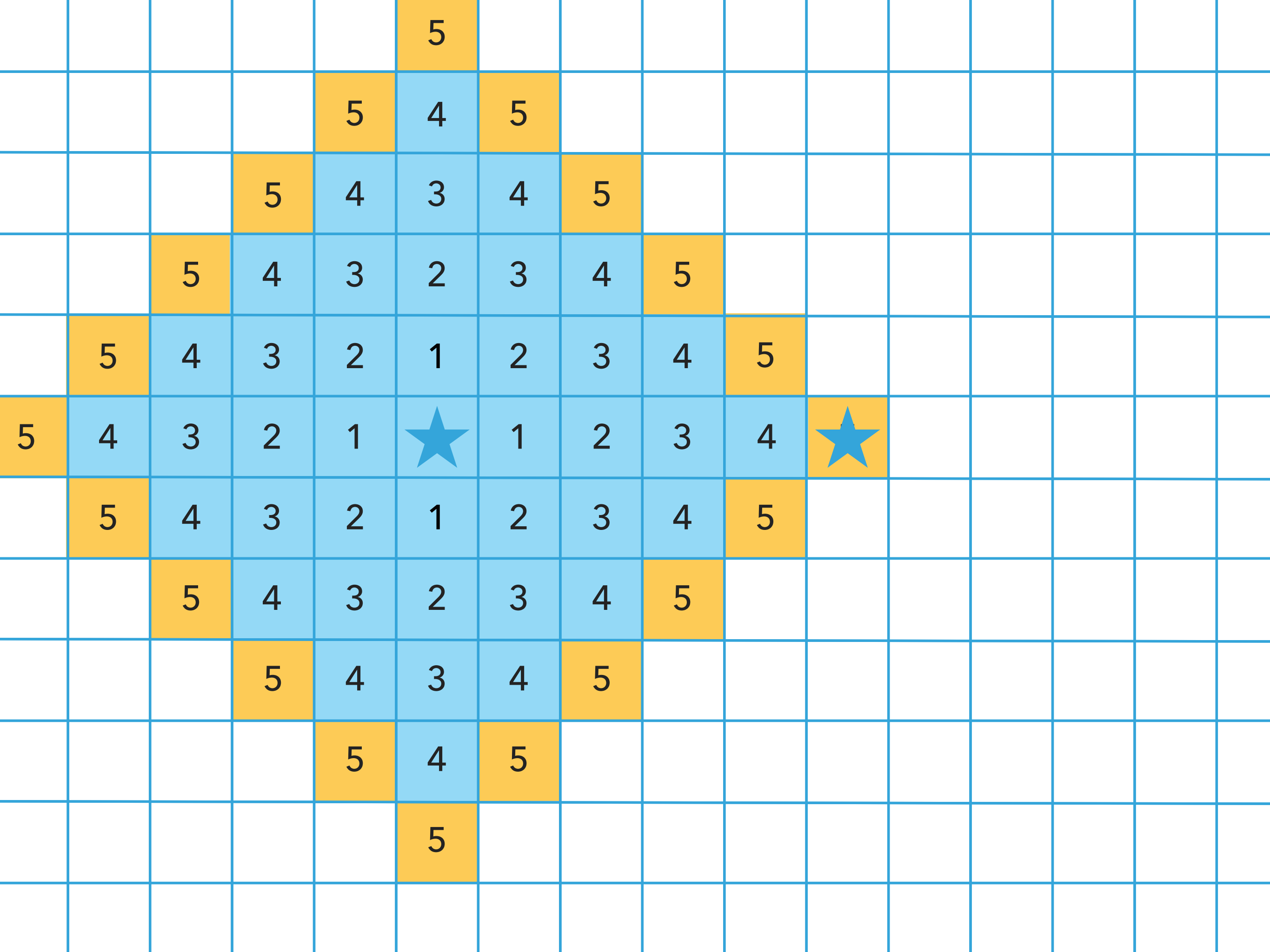


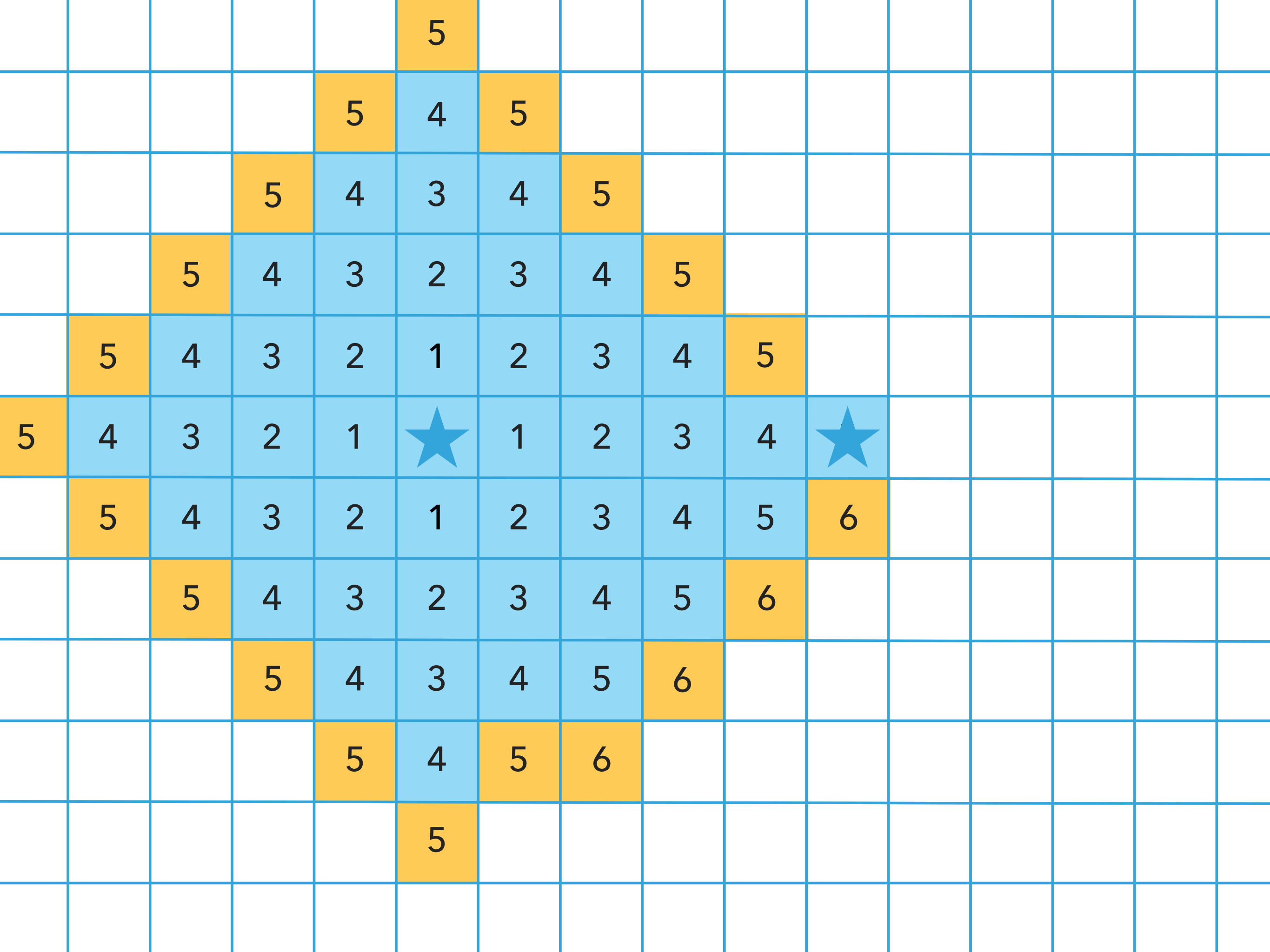










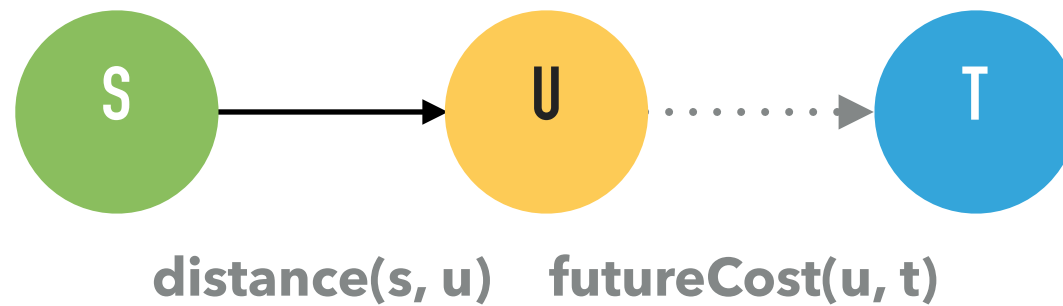


**DIJKSTRA'S MEASURES THE DISTANCE FROM
THE START NODE TO THE CURRENT NODE.**

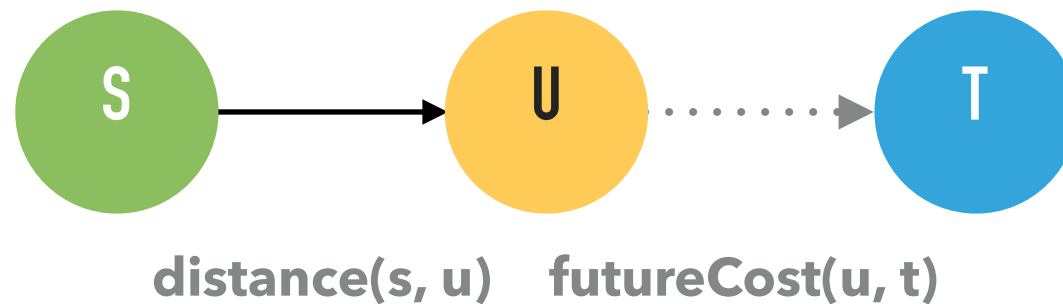
**WE WANT THE DISTANCE FROM THE CURRENT
NODE TO THE DESTINATION.**

SEEING THE
FUTURE

FORMAL DEFINITIONS



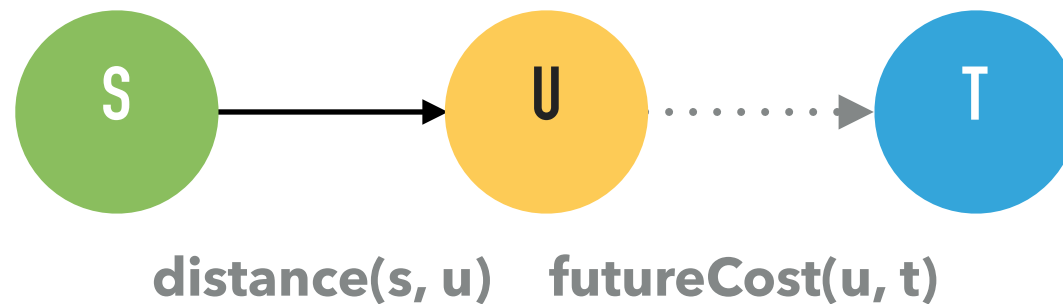
FORMAL DEFINITIONS



DIJKSTRA'S

$$\text{priority}(u) = \text{distance}(s, u)$$

FORMAL DEFINITIONS

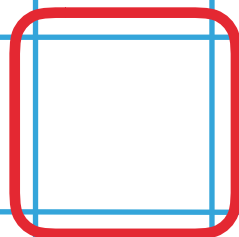


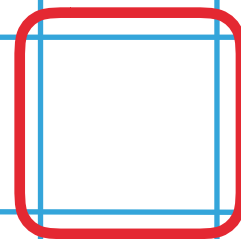
DIJKSTRA'S

$$\text{priority}(u) = \text{distance}(s, u)$$

IDEAL

$$\text{priority}(u) = \text{distance}(s, u) + \text{futureCost}(u, t)$$





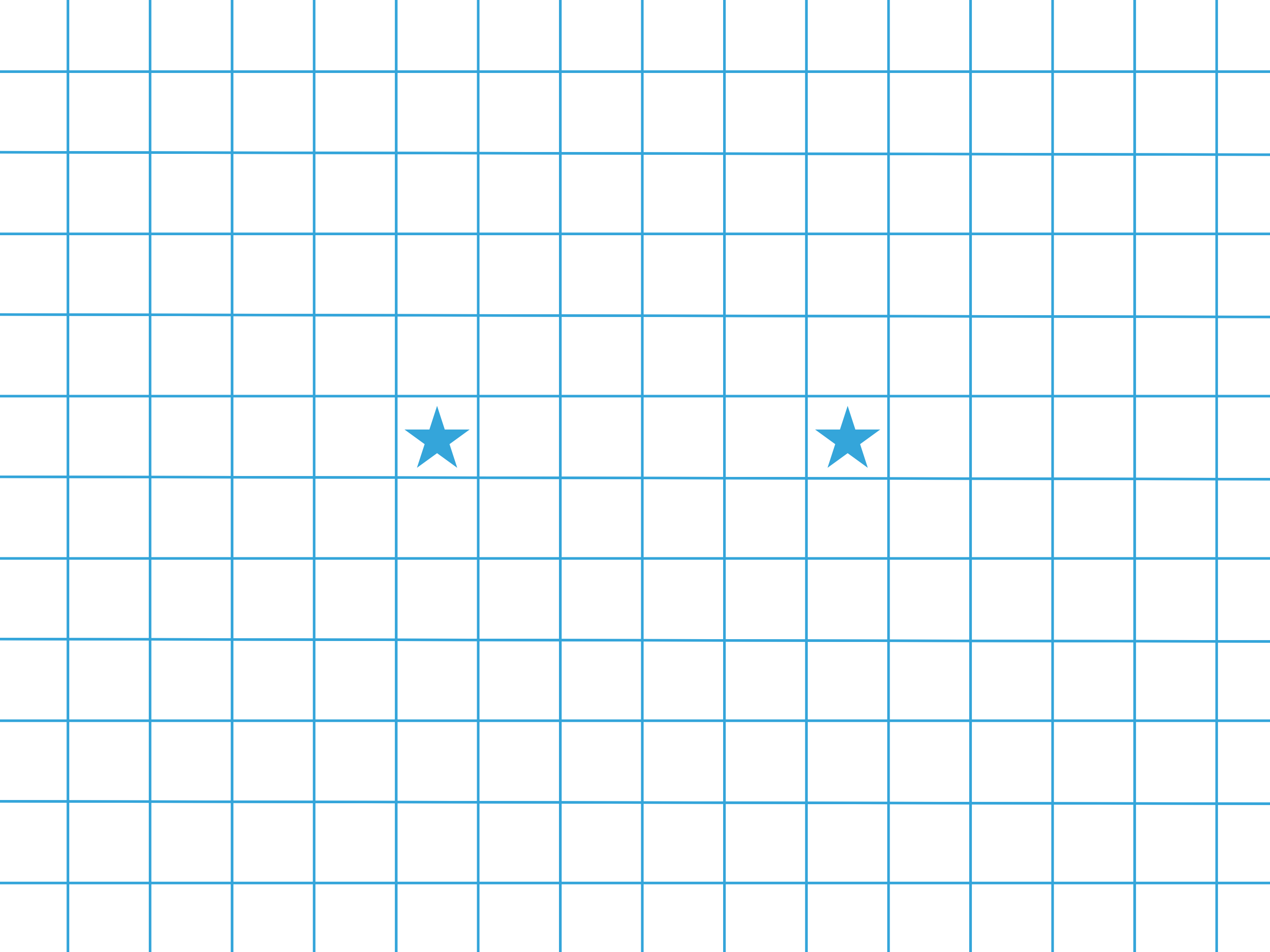
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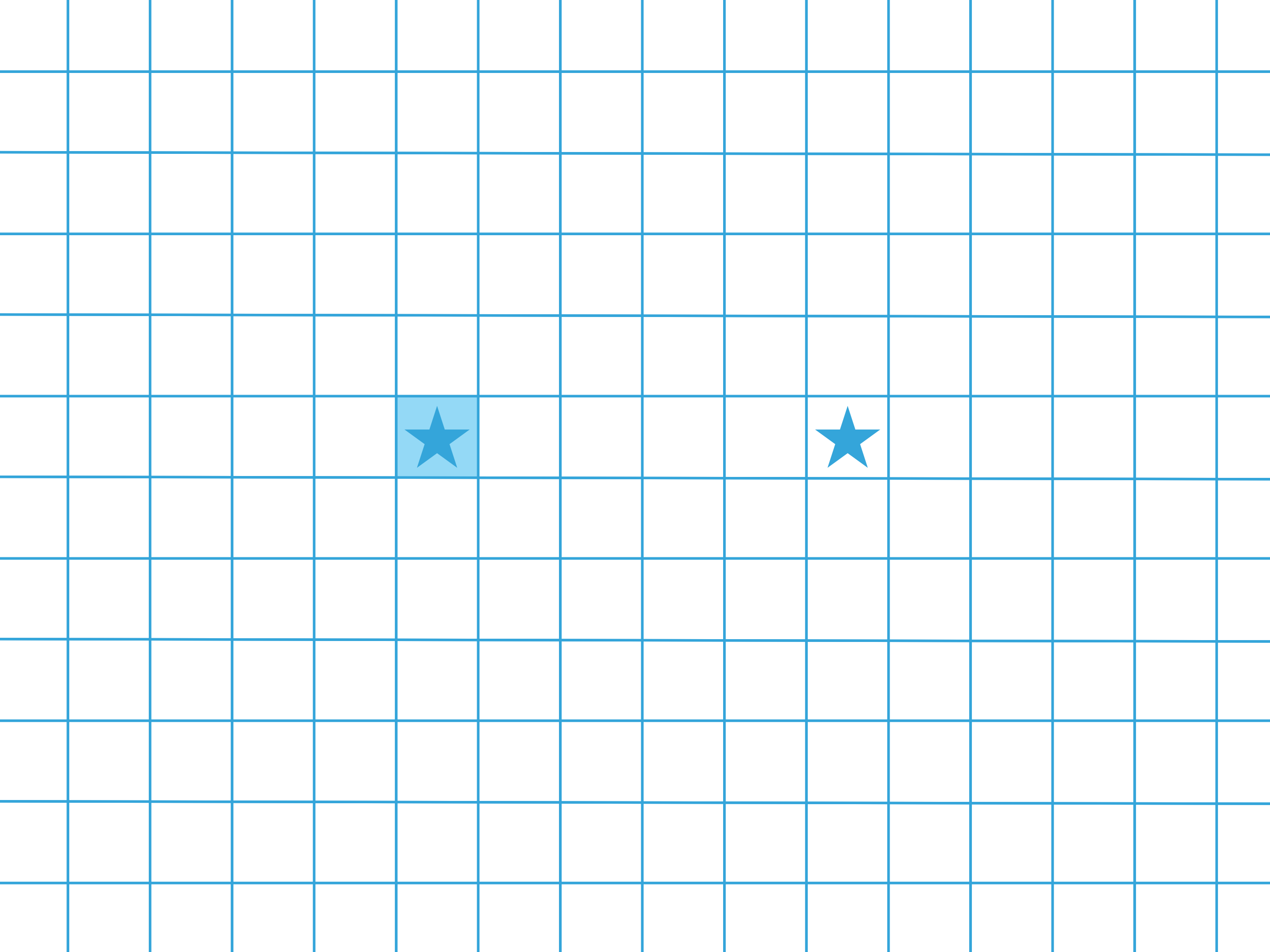


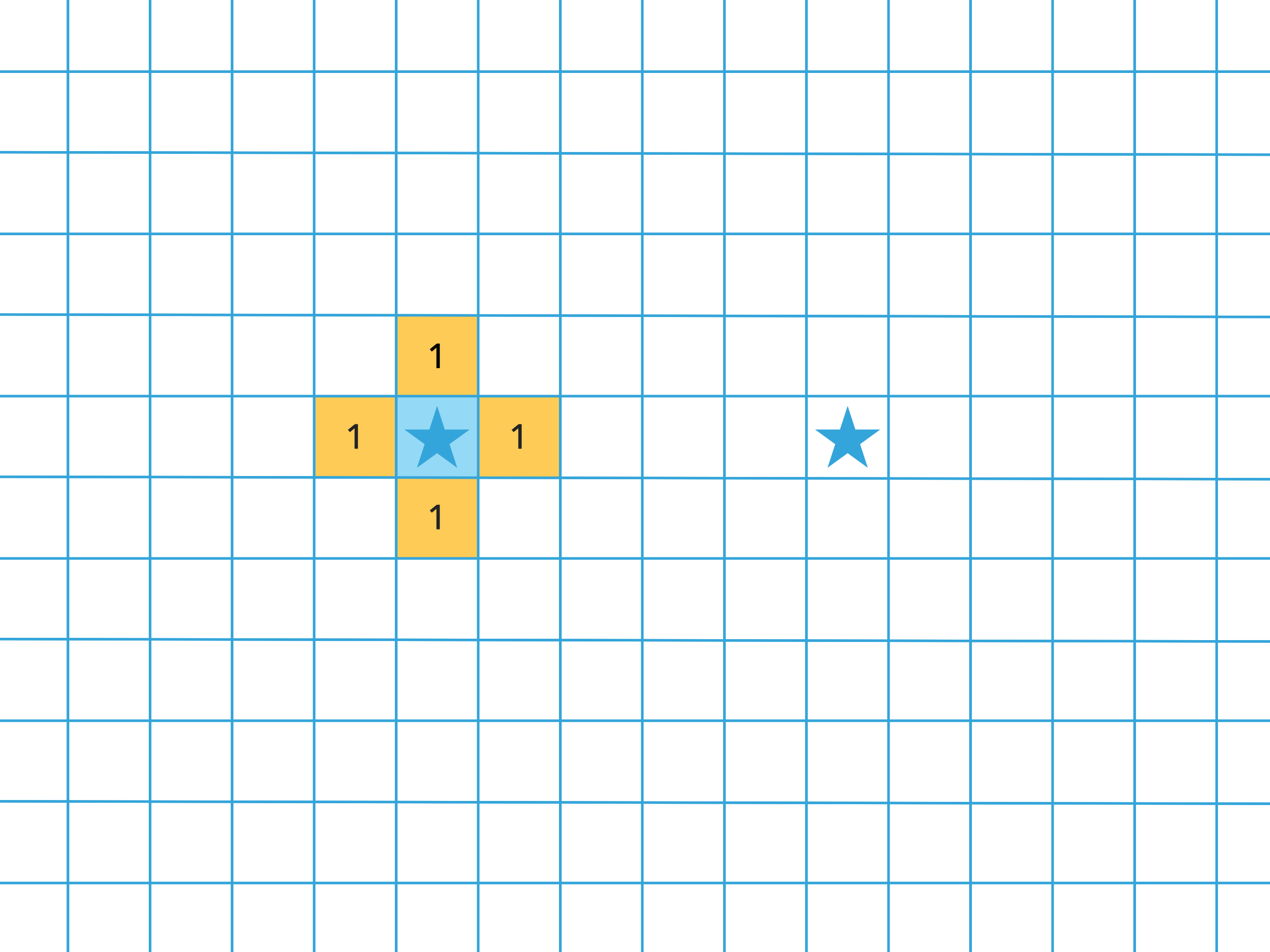
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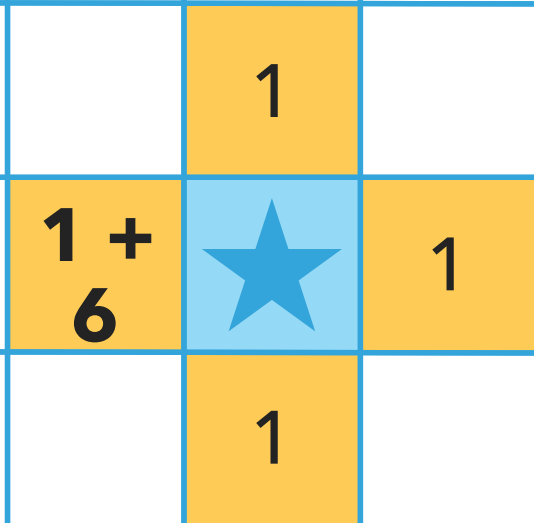


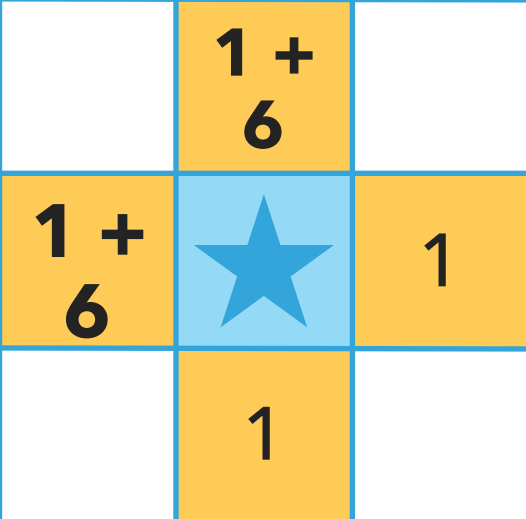
```
function futureCost(u, t)
  return abs(u.row - t.row) + abs(u.col - t.col)
```





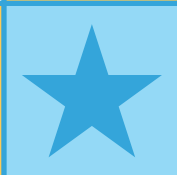






$$\begin{array}{c} 1 + \\ 6 \end{array}$$

$$\begin{array}{c} 1 + \\ 6 \end{array}$$



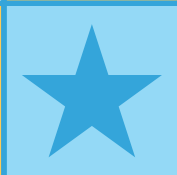
$$1$$

$$\begin{array}{c} 1 + \\ 6 \end{array}$$



$$\begin{array}{r} 1 + \\ 6 \end{array}$$

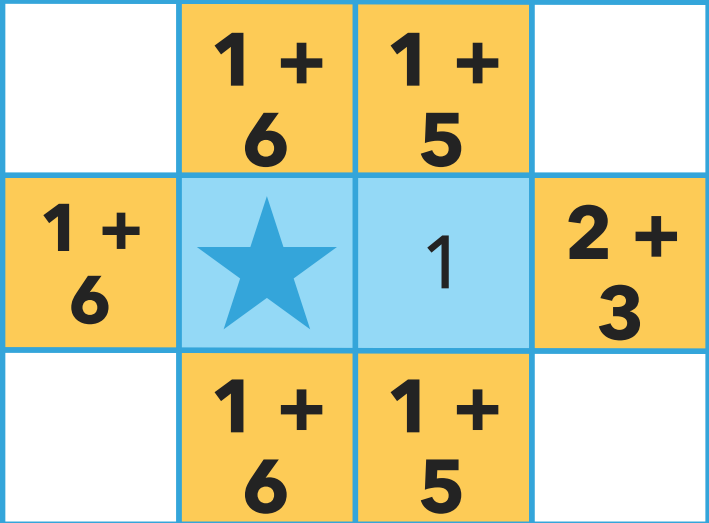
$$\begin{array}{r} 1 + \\ 6 \end{array}$$

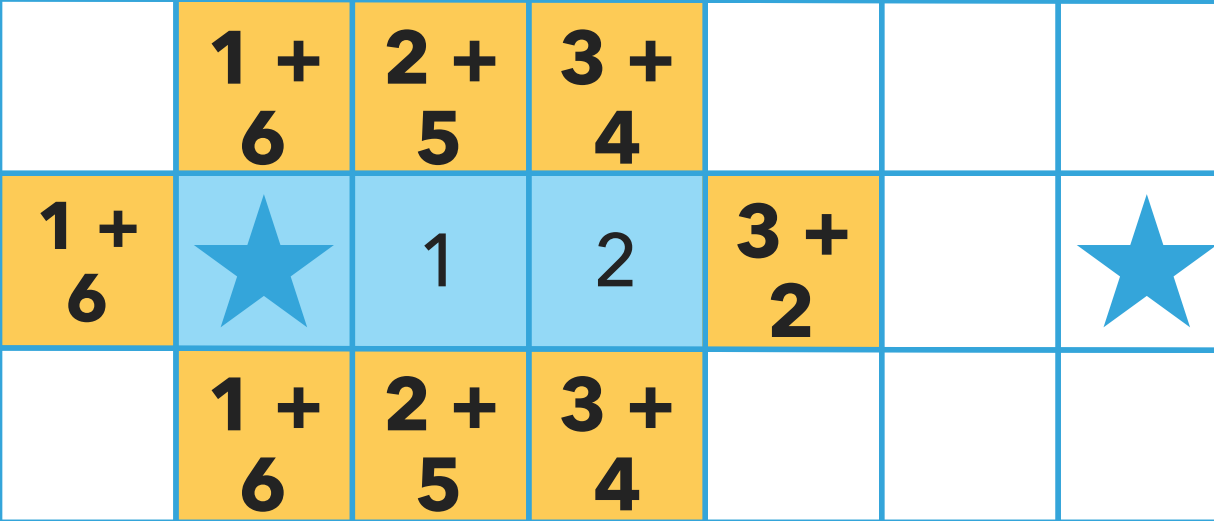


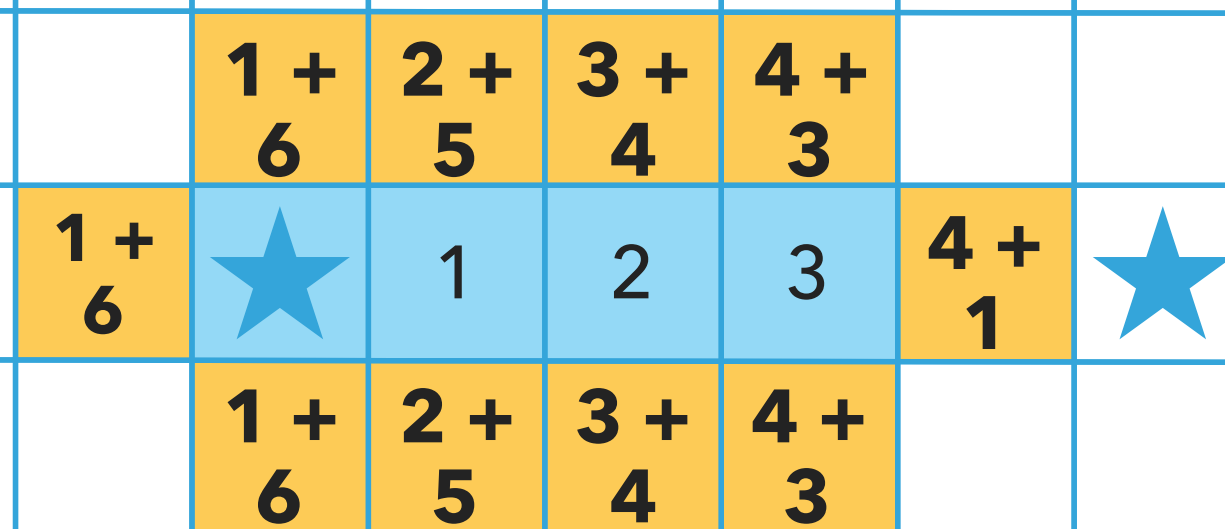
$$\begin{array}{r} 1 + \\ 4 \end{array}$$

$$\begin{array}{r} 1 + \\ 6 \end{array}$$



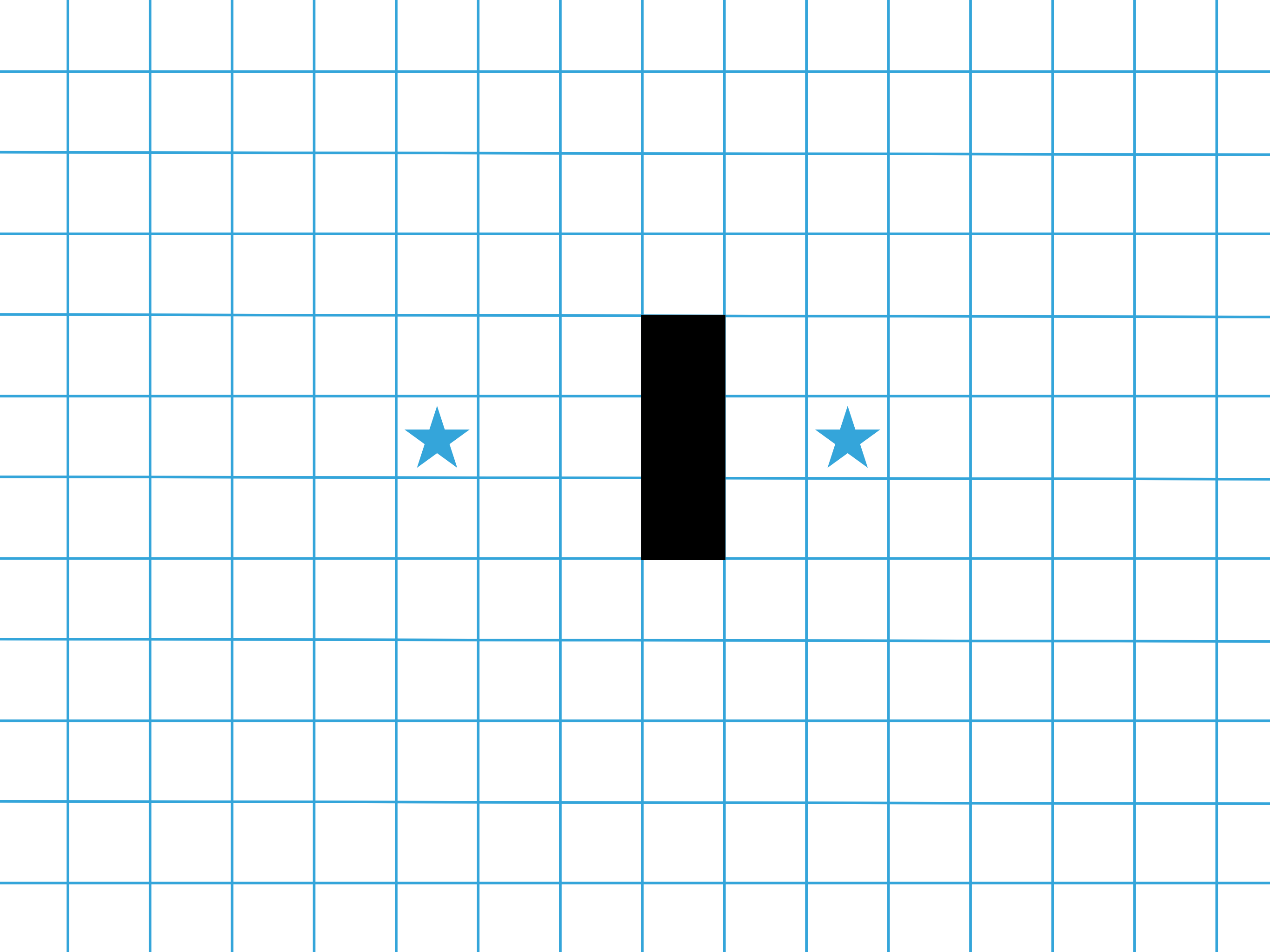






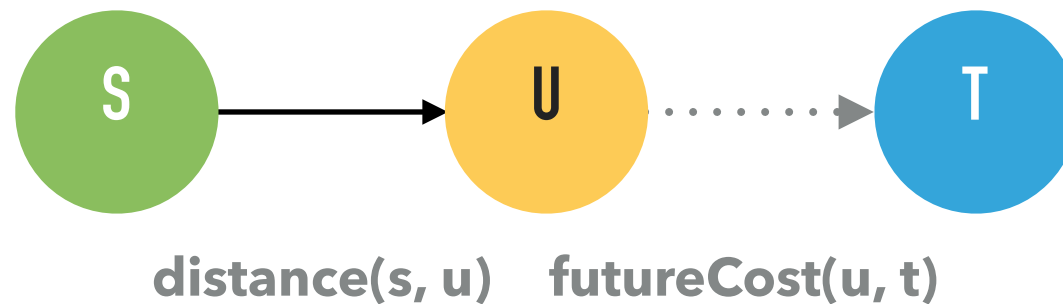
[illegible]

					1 + 6	2 + 5	3 + 4	4 + 3	5 + 2					
				1 + 6	★	1	2	3	4	★				
					1 + 6	2 + 5	3 + 4	4 + 3	5 + 2					



**MAKING GOOD
LIFE DECISIONS**

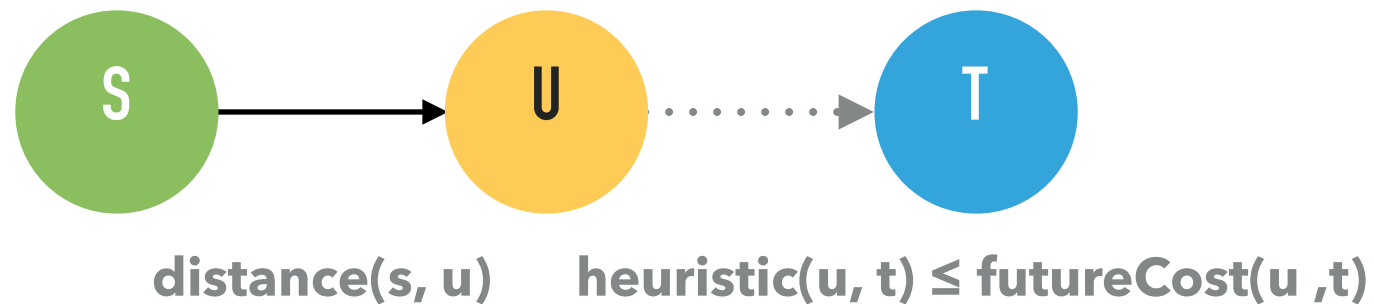
FORMAL DEFINITIONS



IDEAL

$$\text{priority}(u) = \text{distance}(s, u) + \text{futureCost}(u, t)$$

FORMAL DEFINITIONS



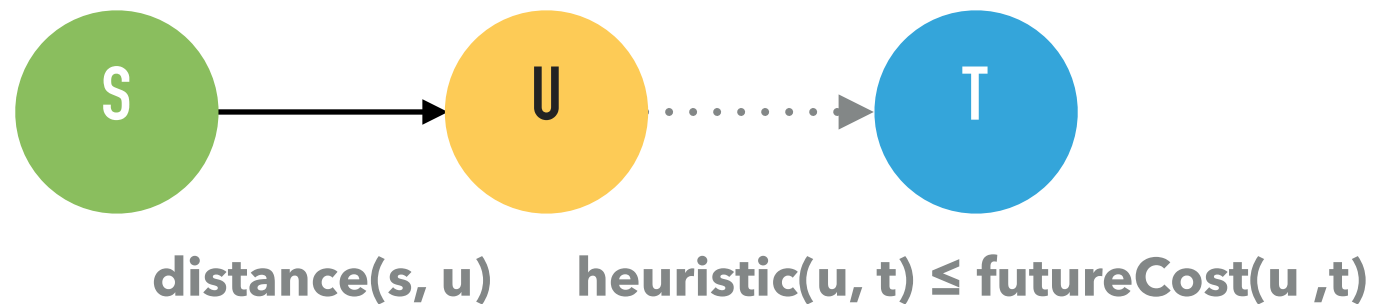
IDEAL

$$\text{priority}(u) = \text{distance}(s, u) + \text{futureCost}(u, t)$$

A*

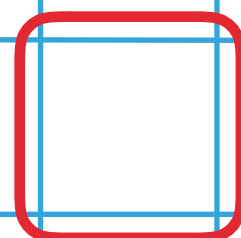
$$\text{priority}(u) = \text{distance}(s, u) + \text{heuristic}(u, t)$$

HEURISTICS



A heuristic is a function that **underestimates** the cost of traveling from u to t.

It's a "relaxation" heuristic.

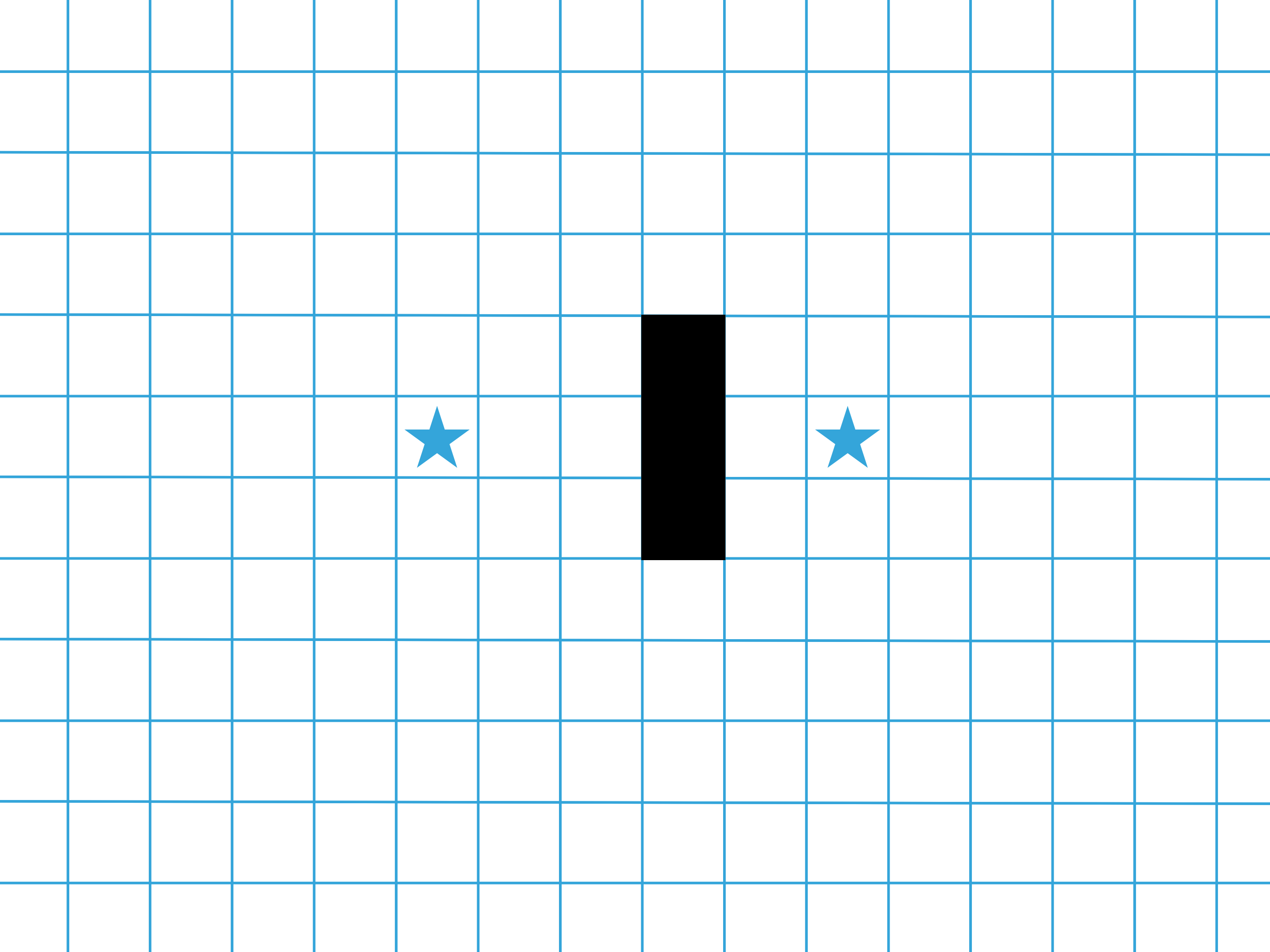


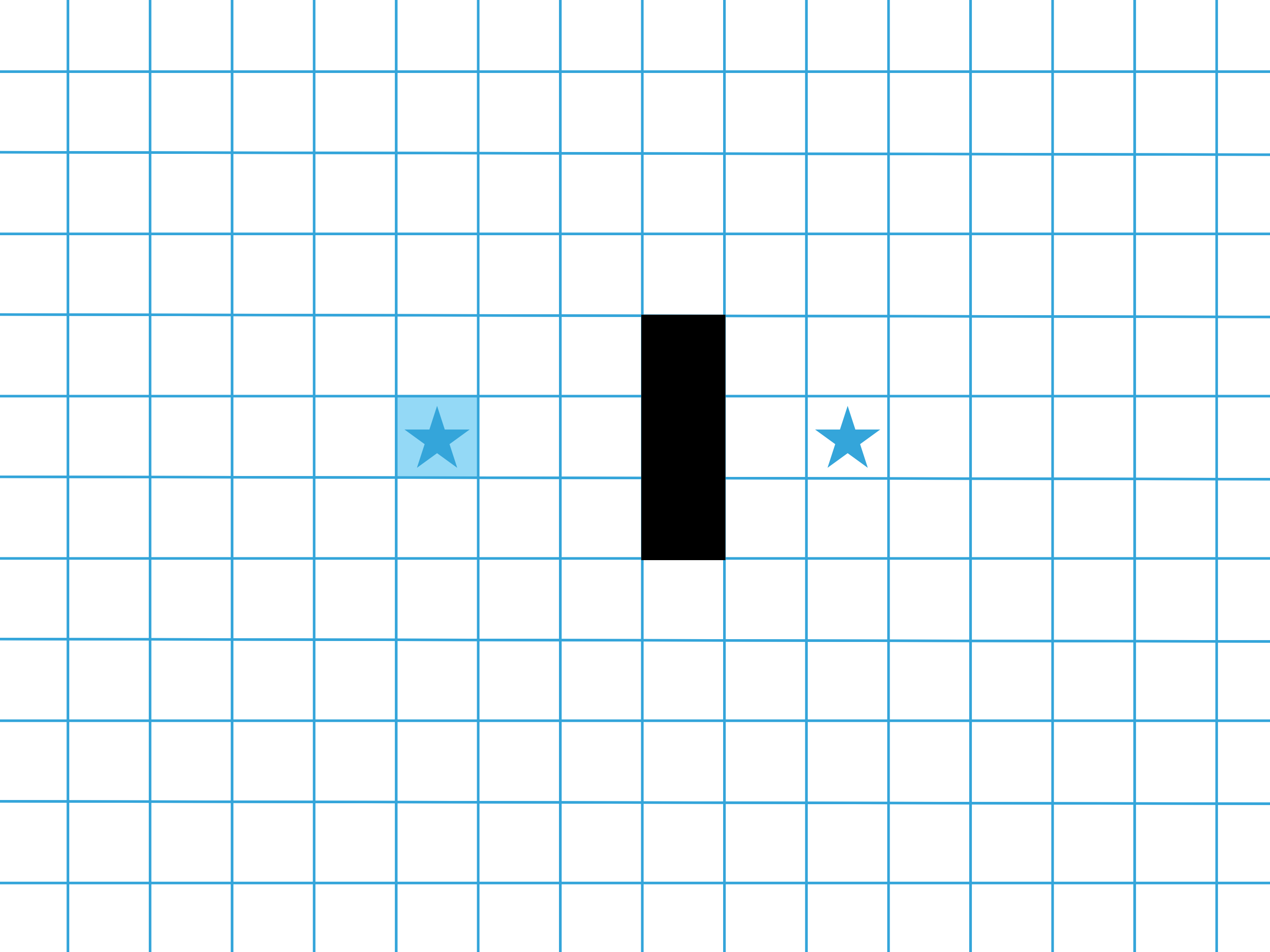
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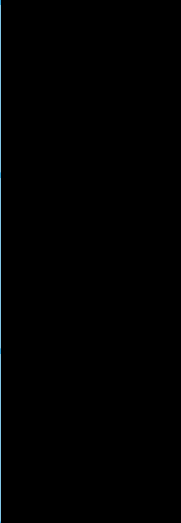
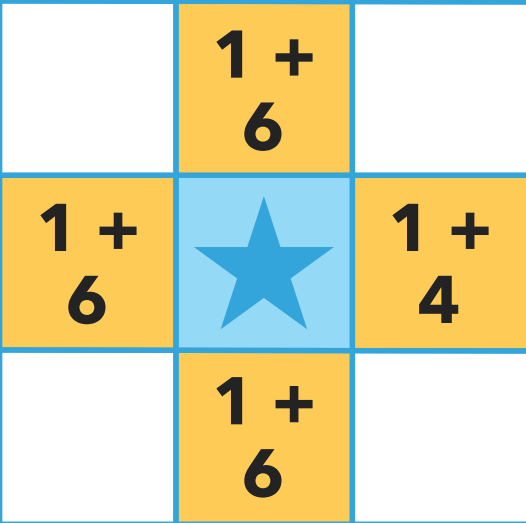


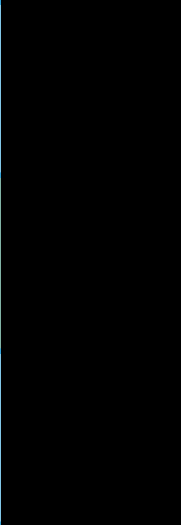
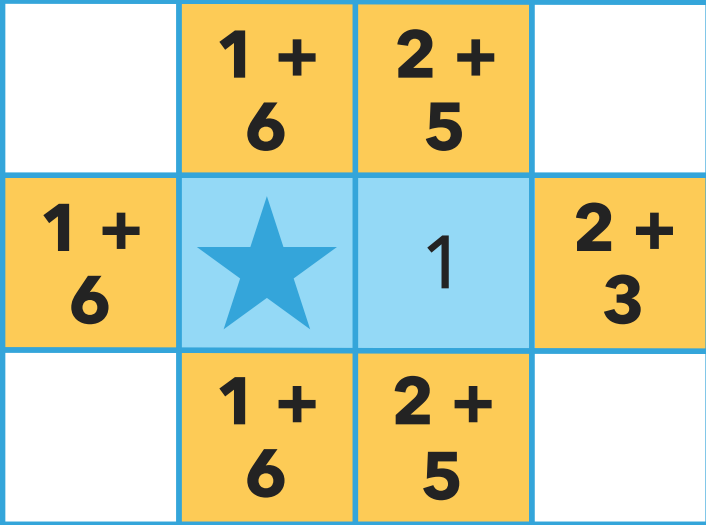
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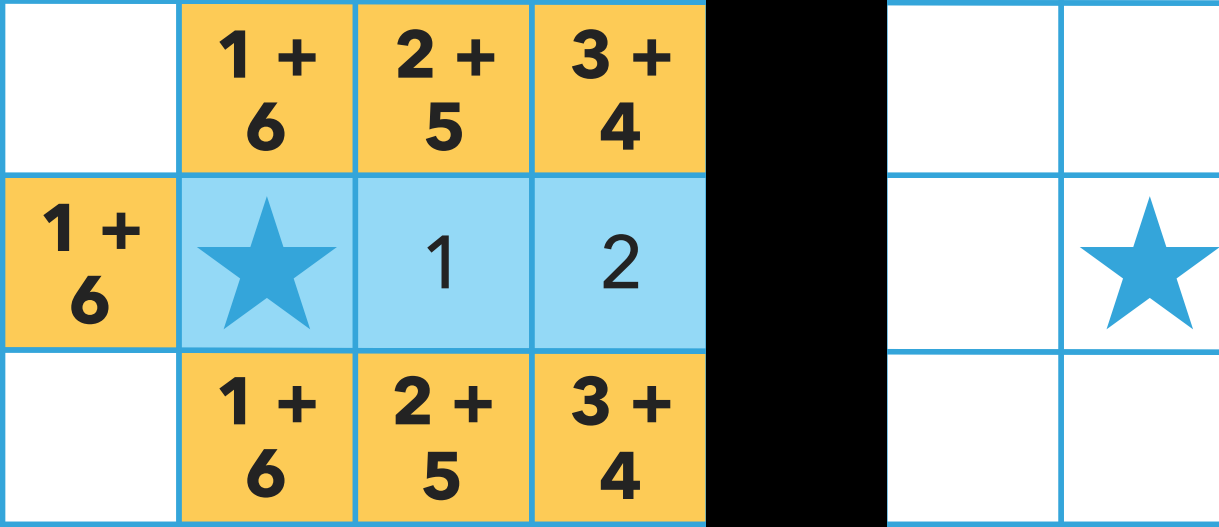


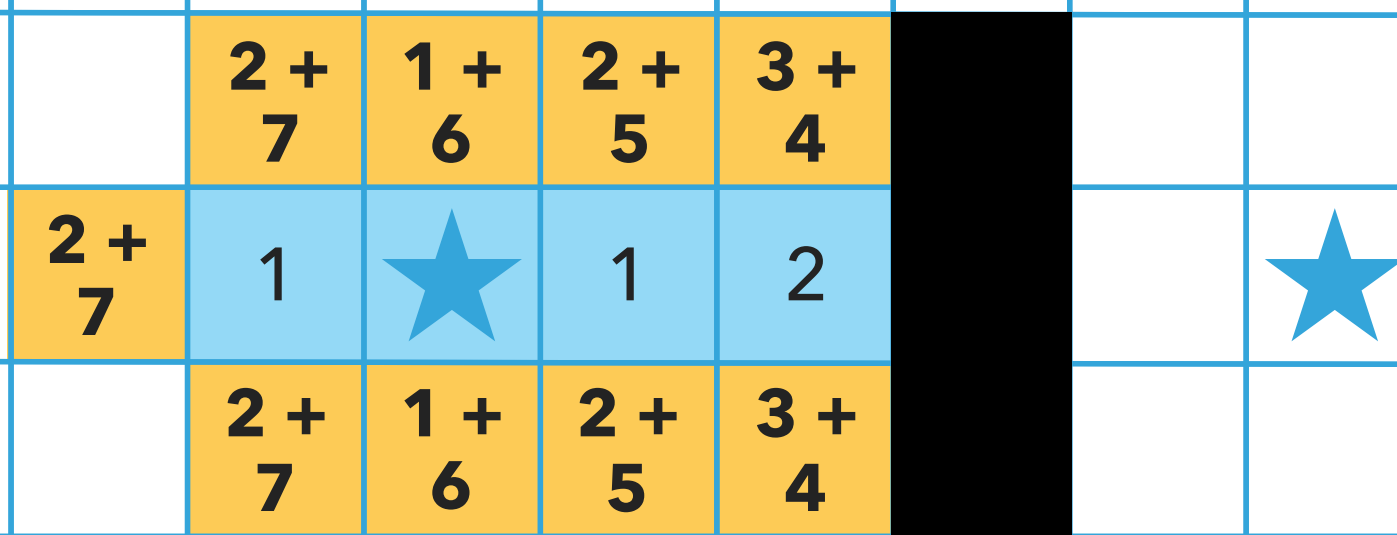


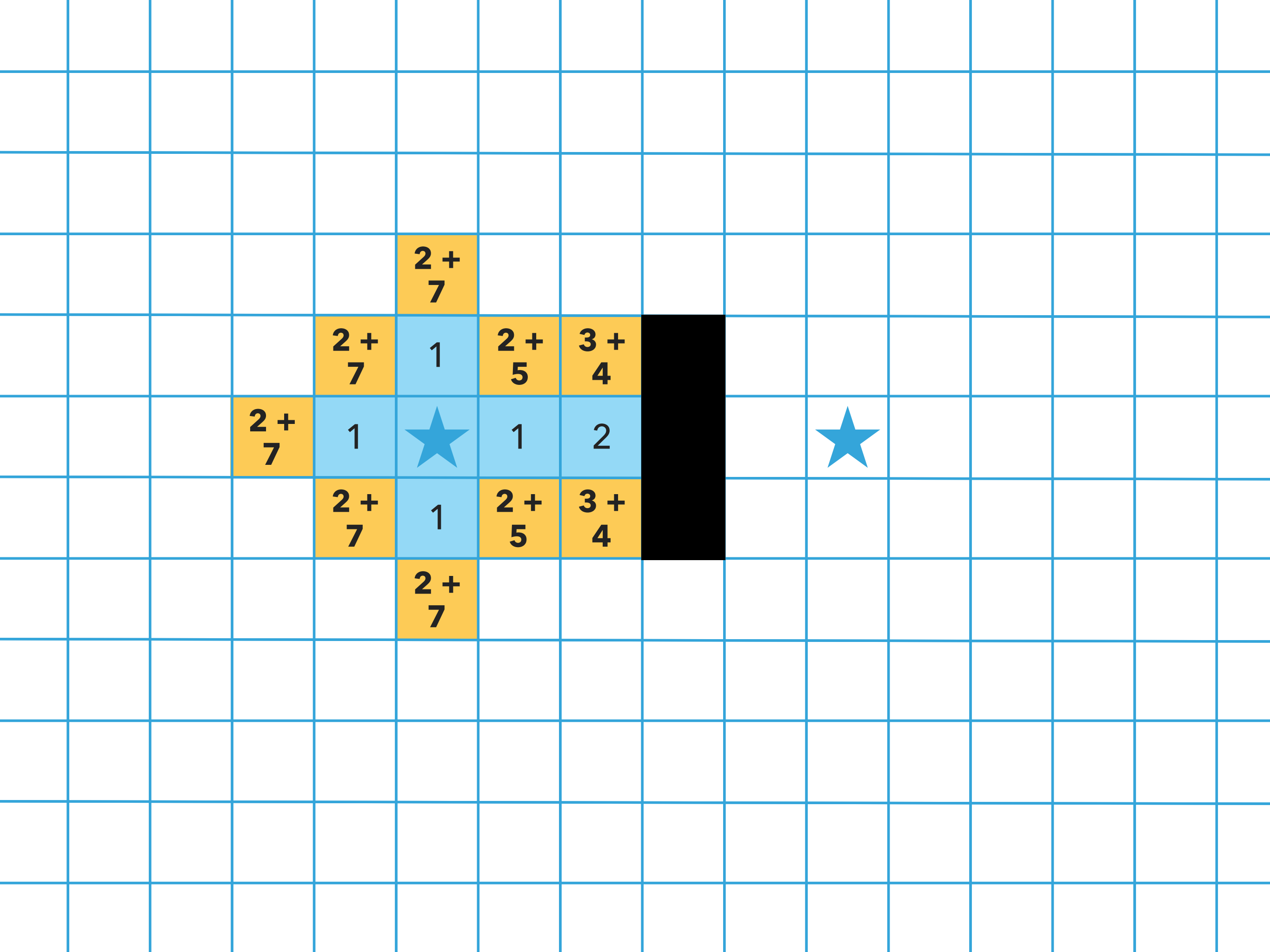


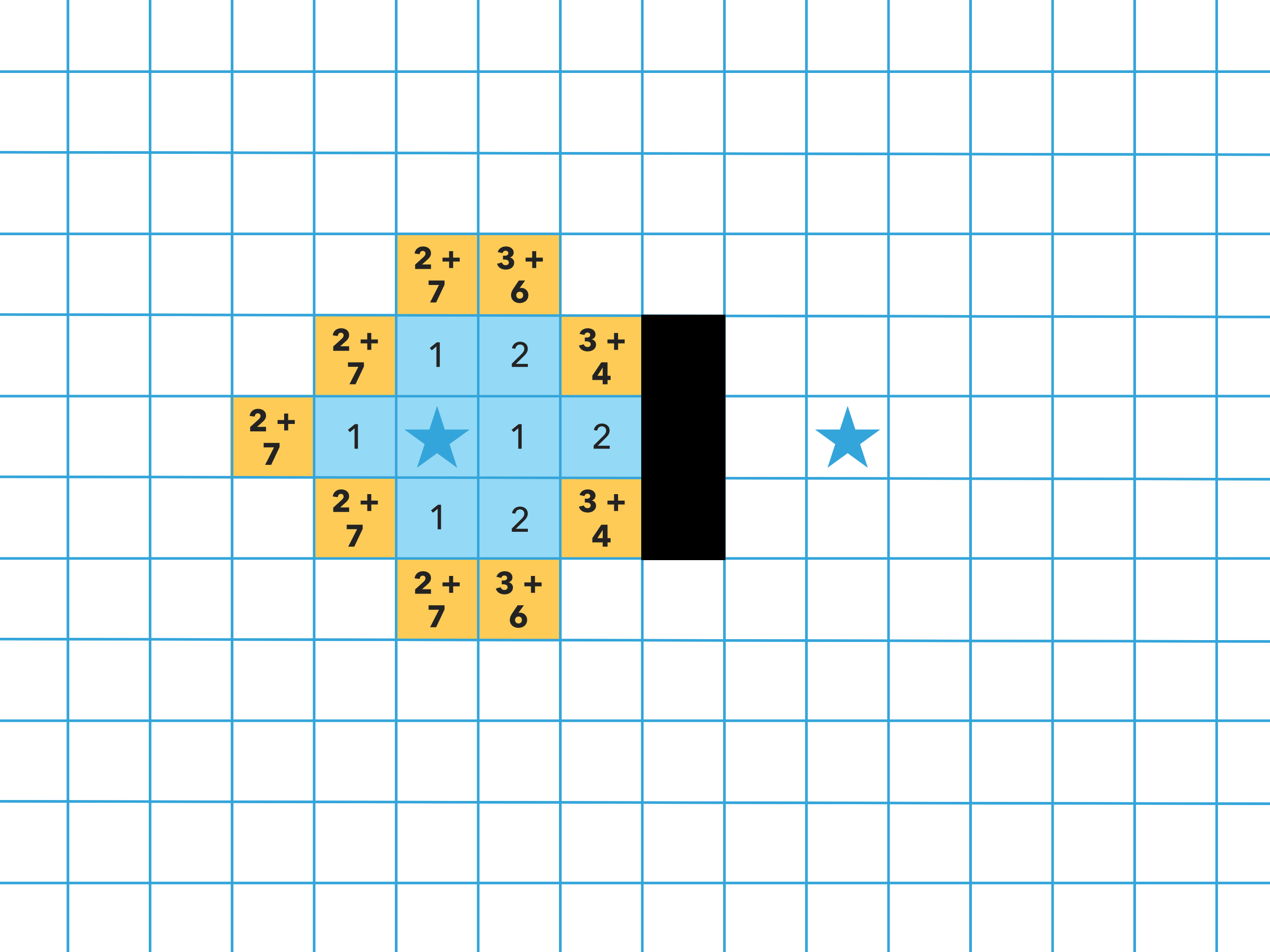


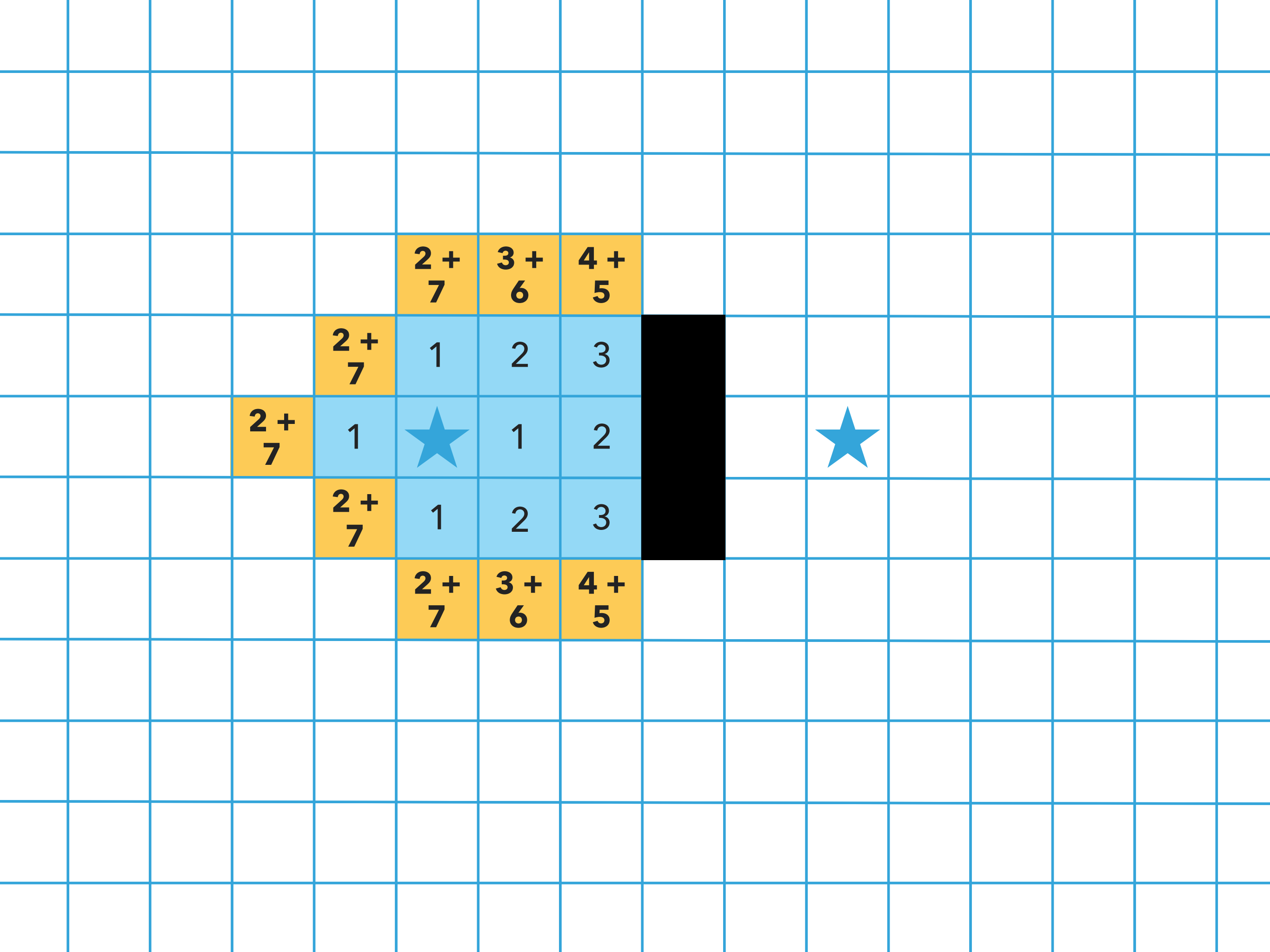


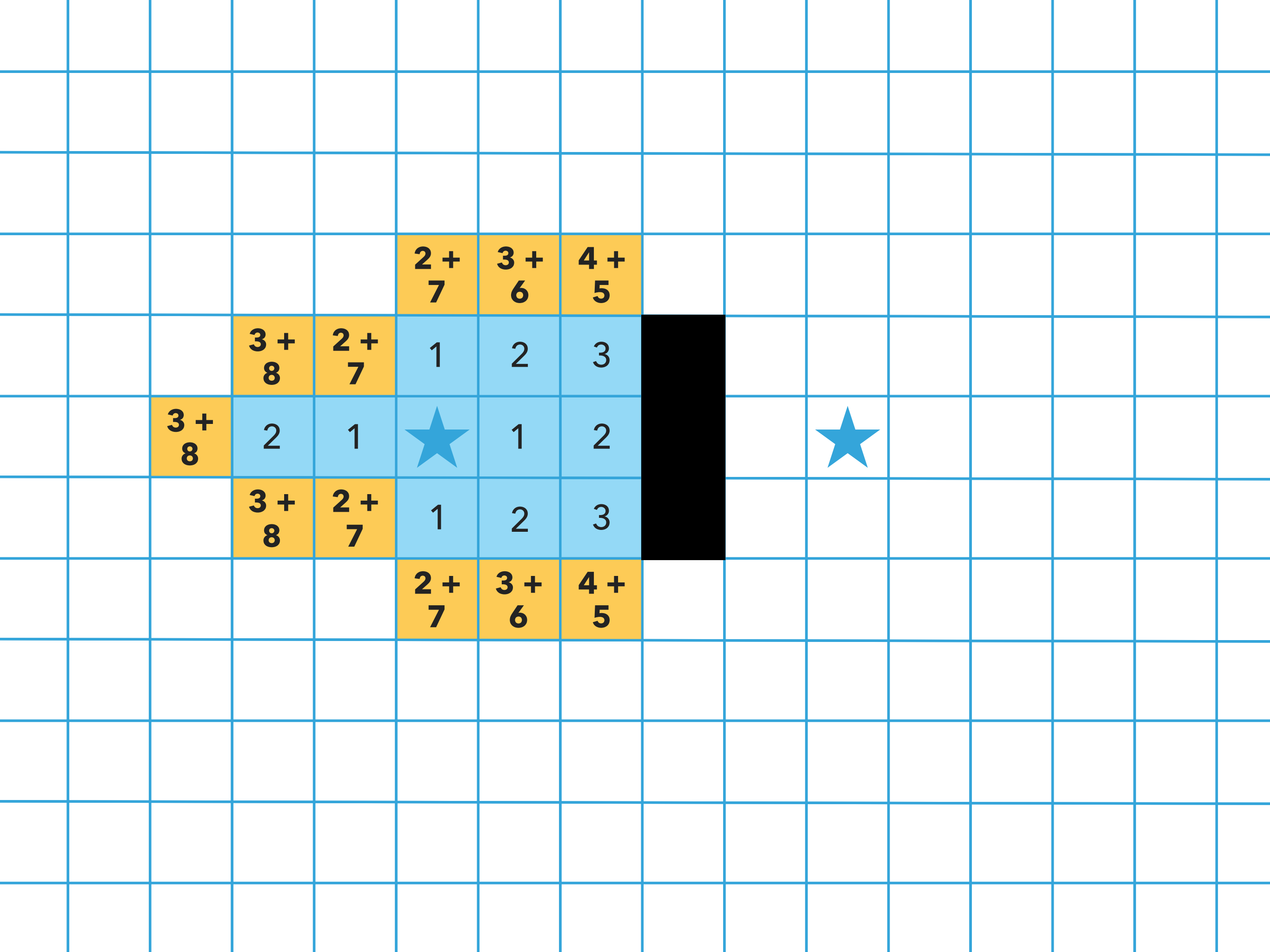


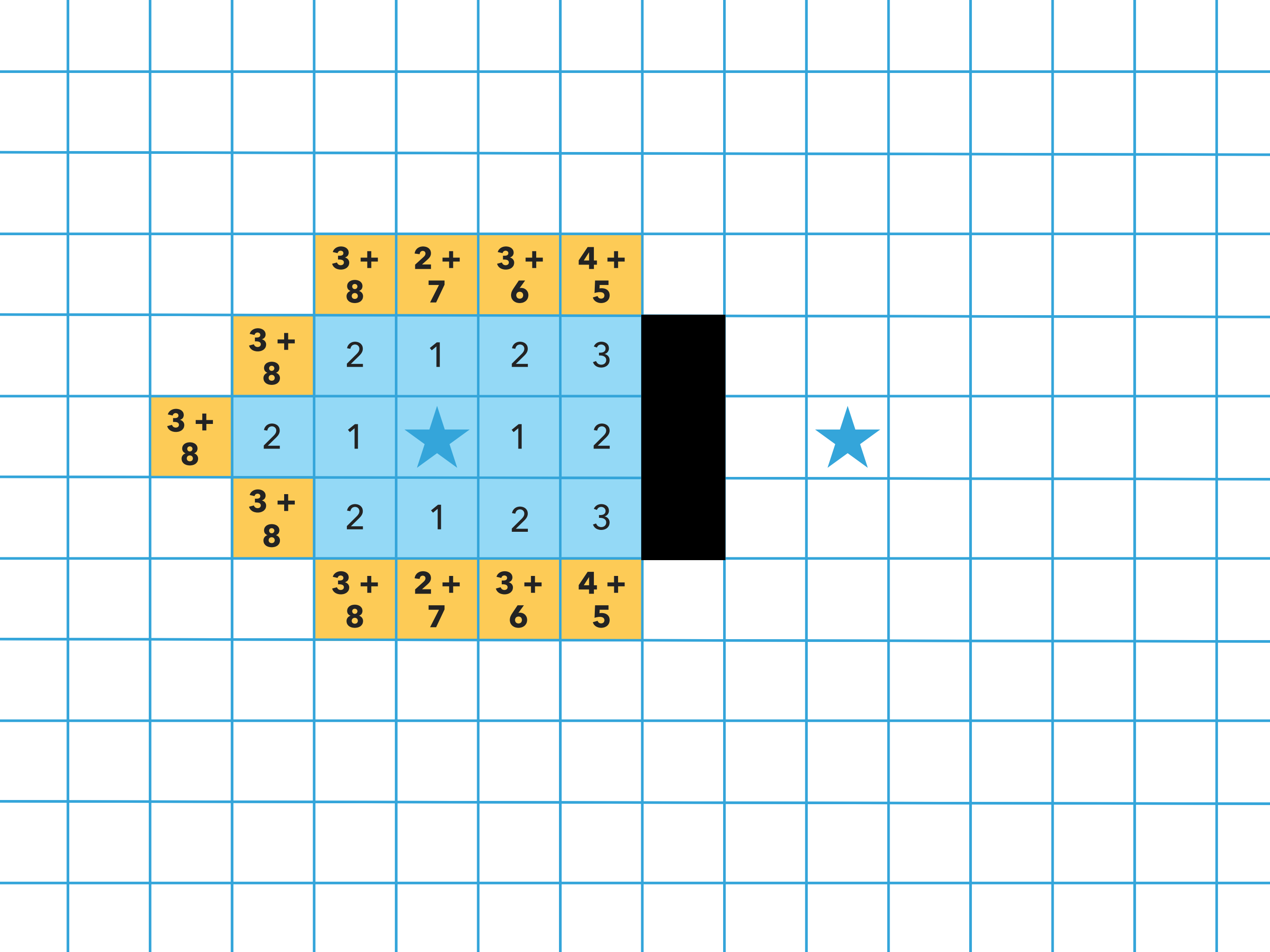


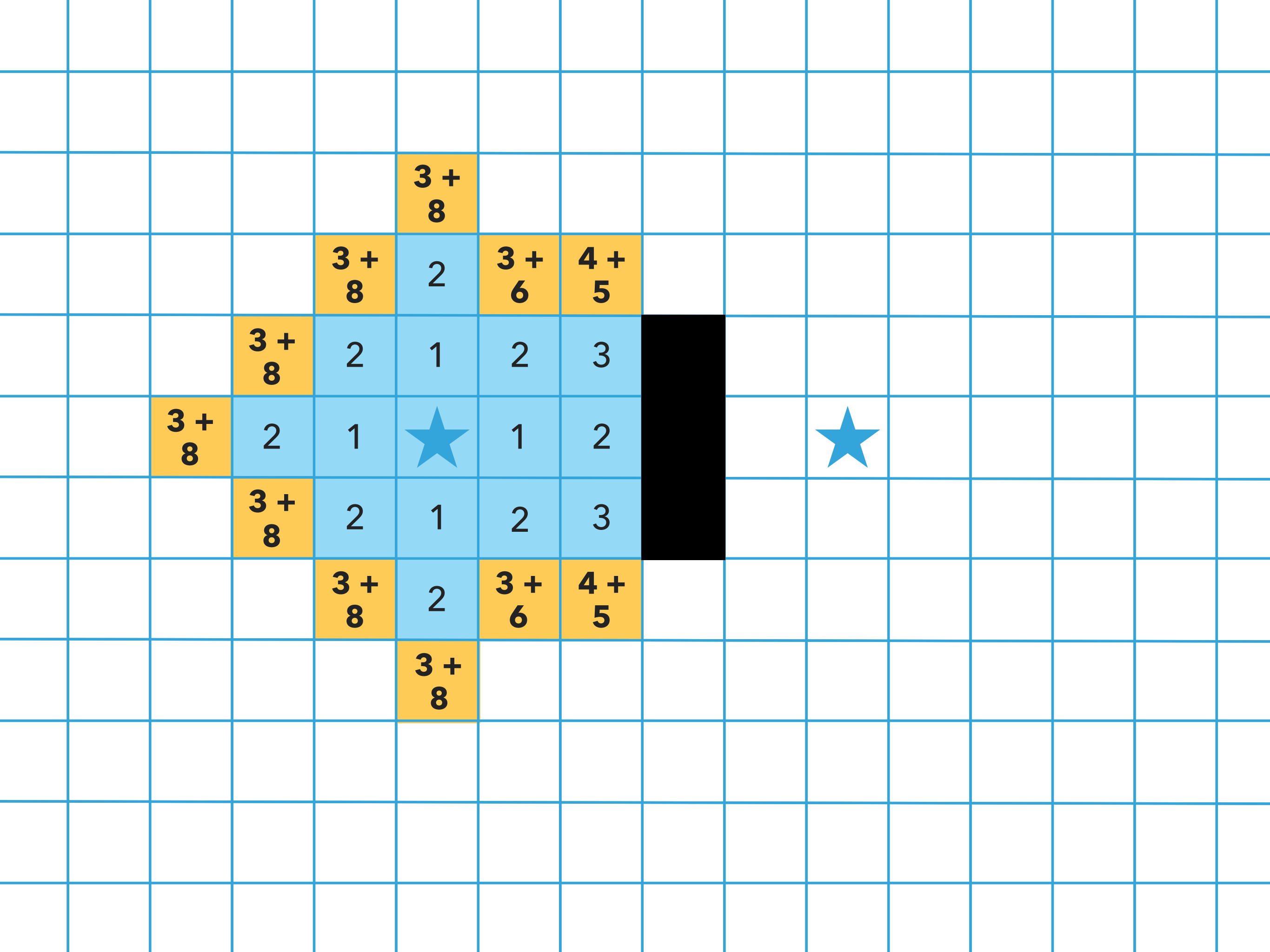


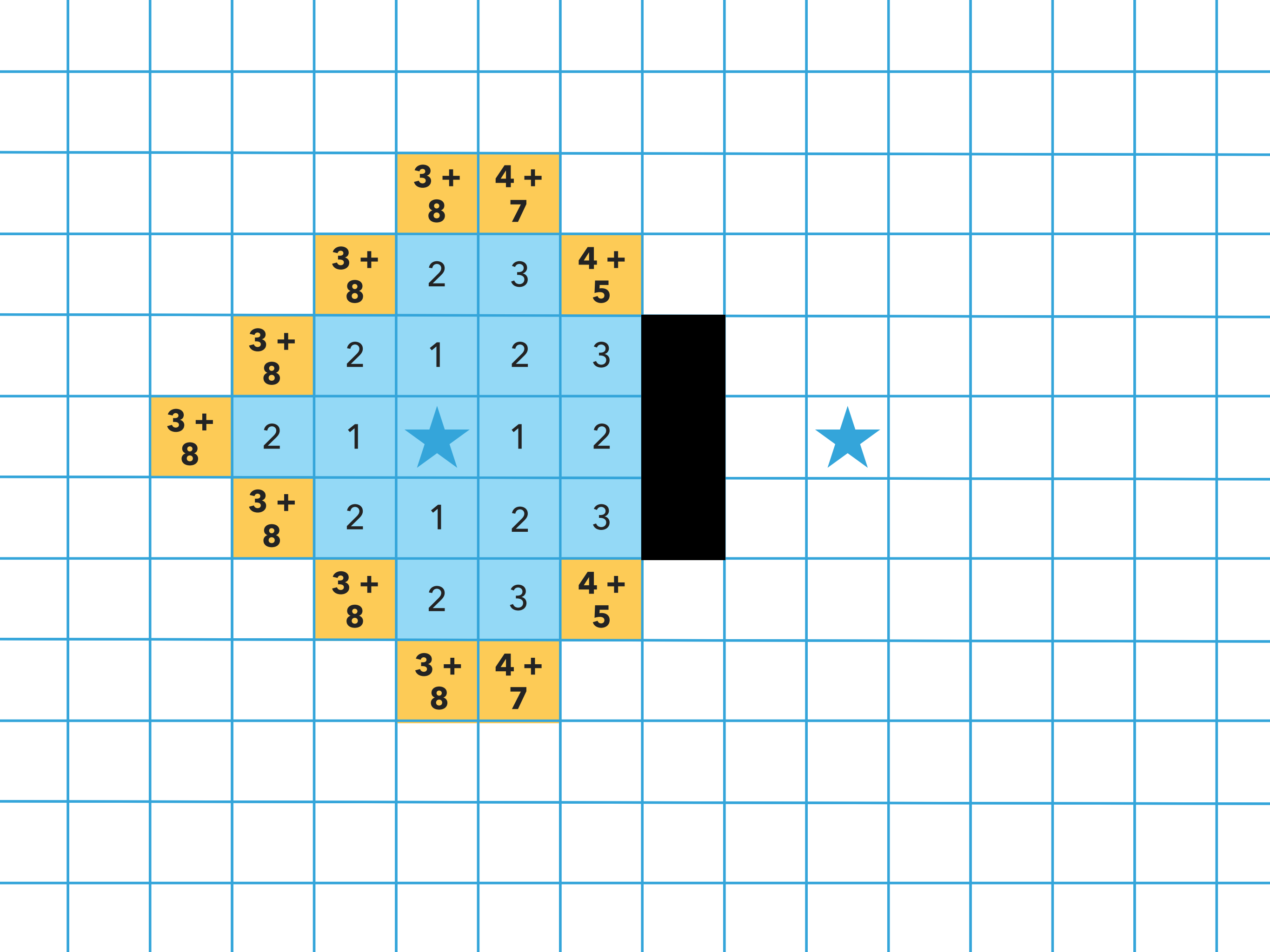


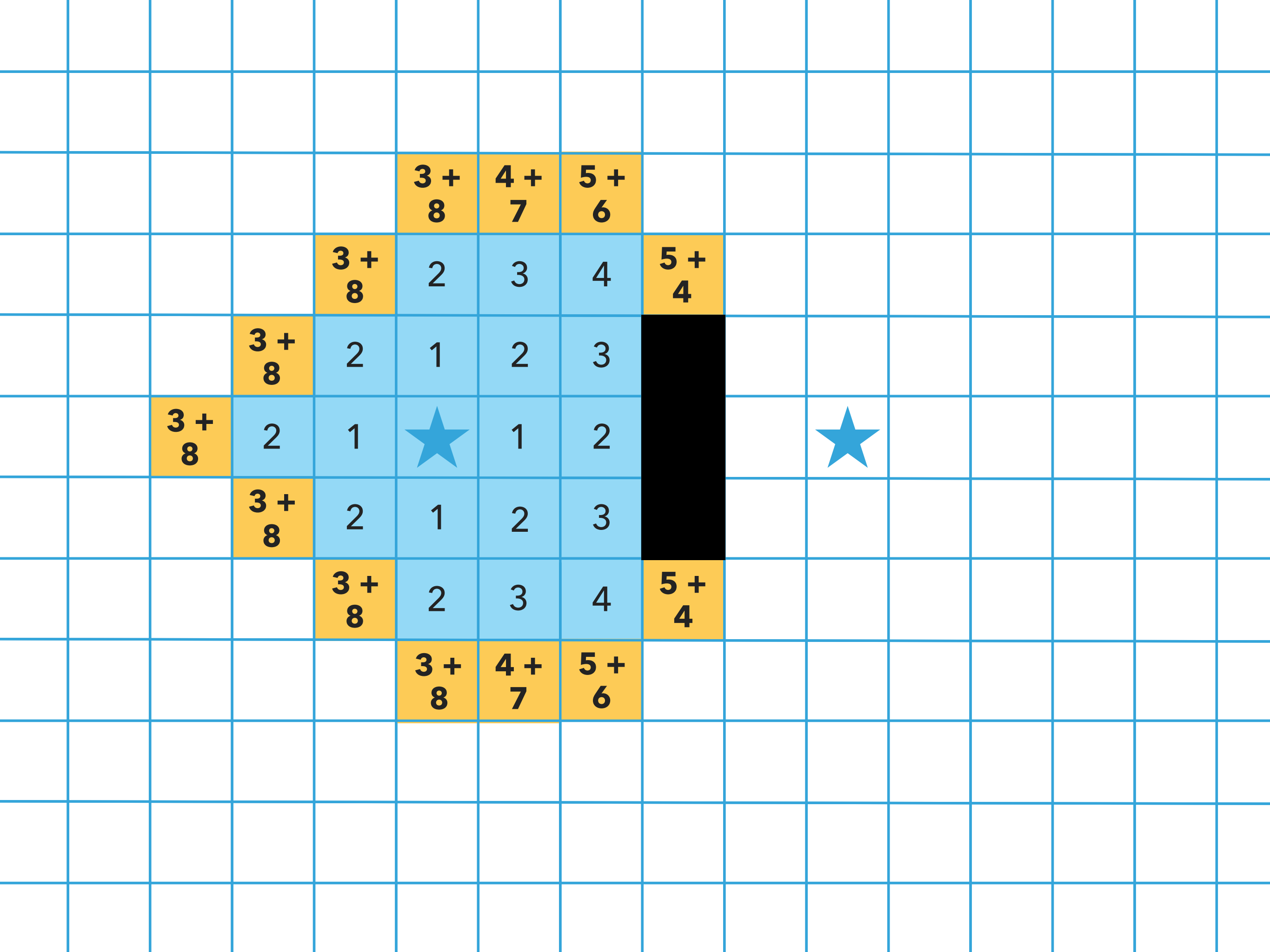


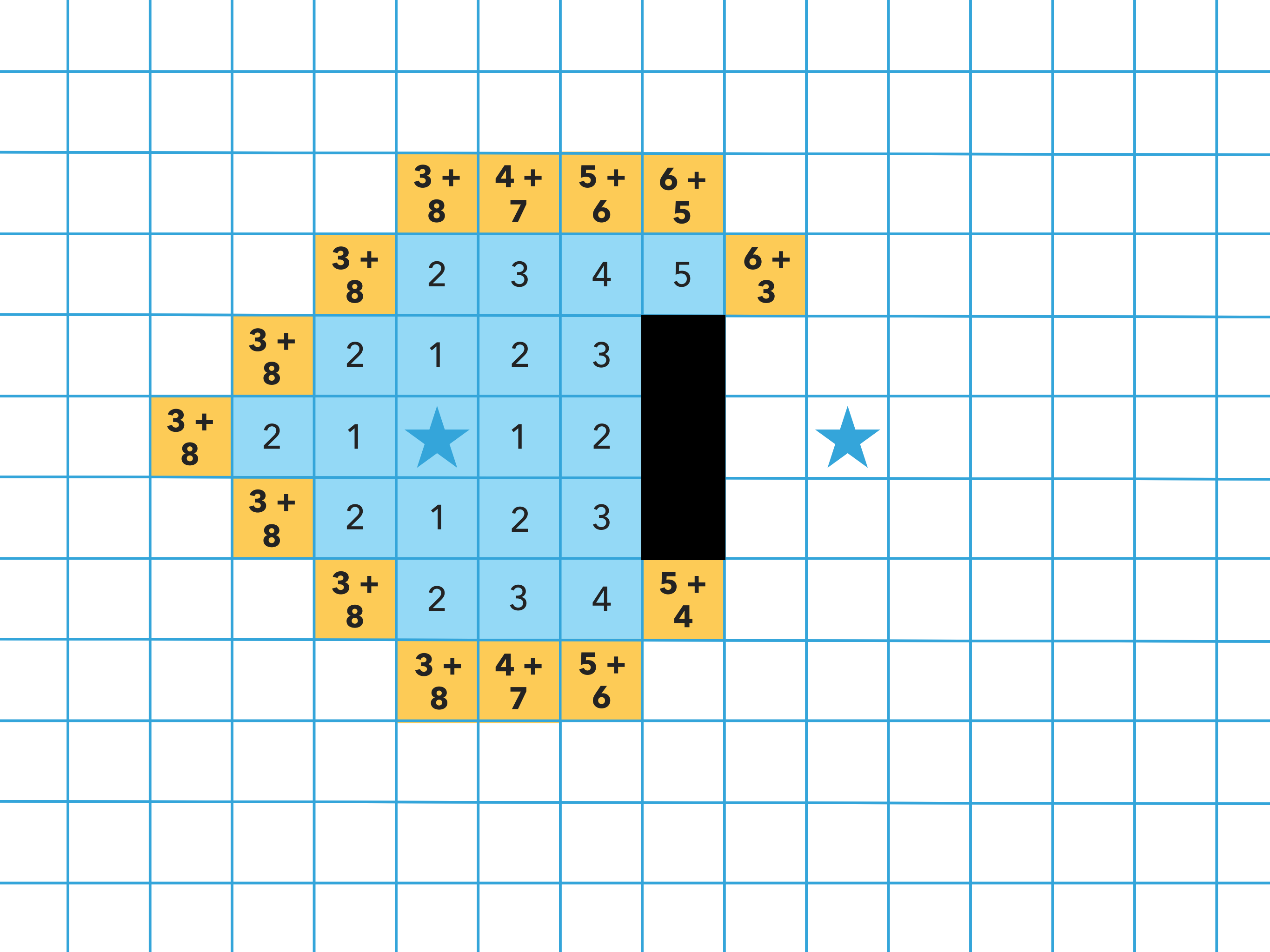


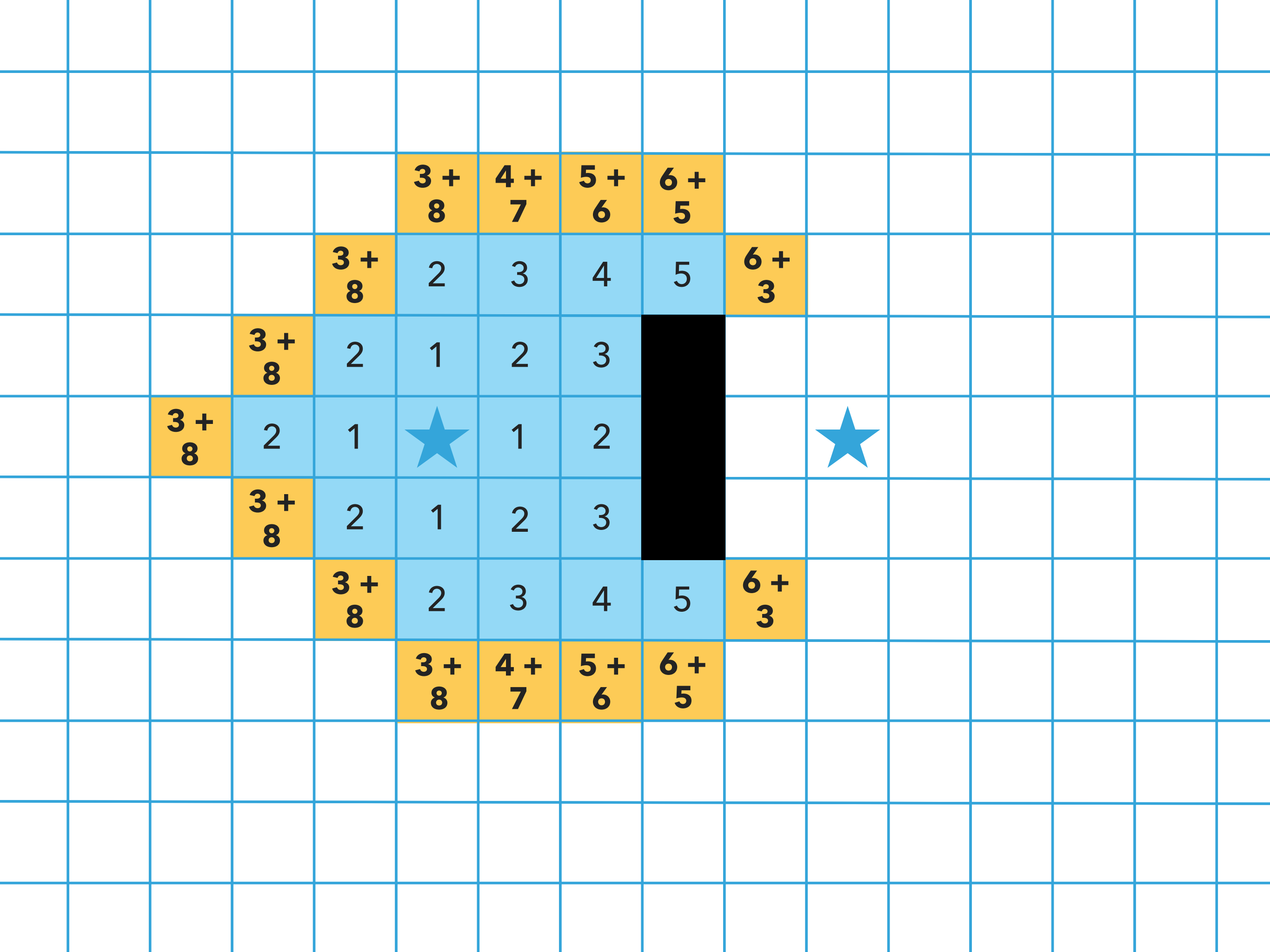


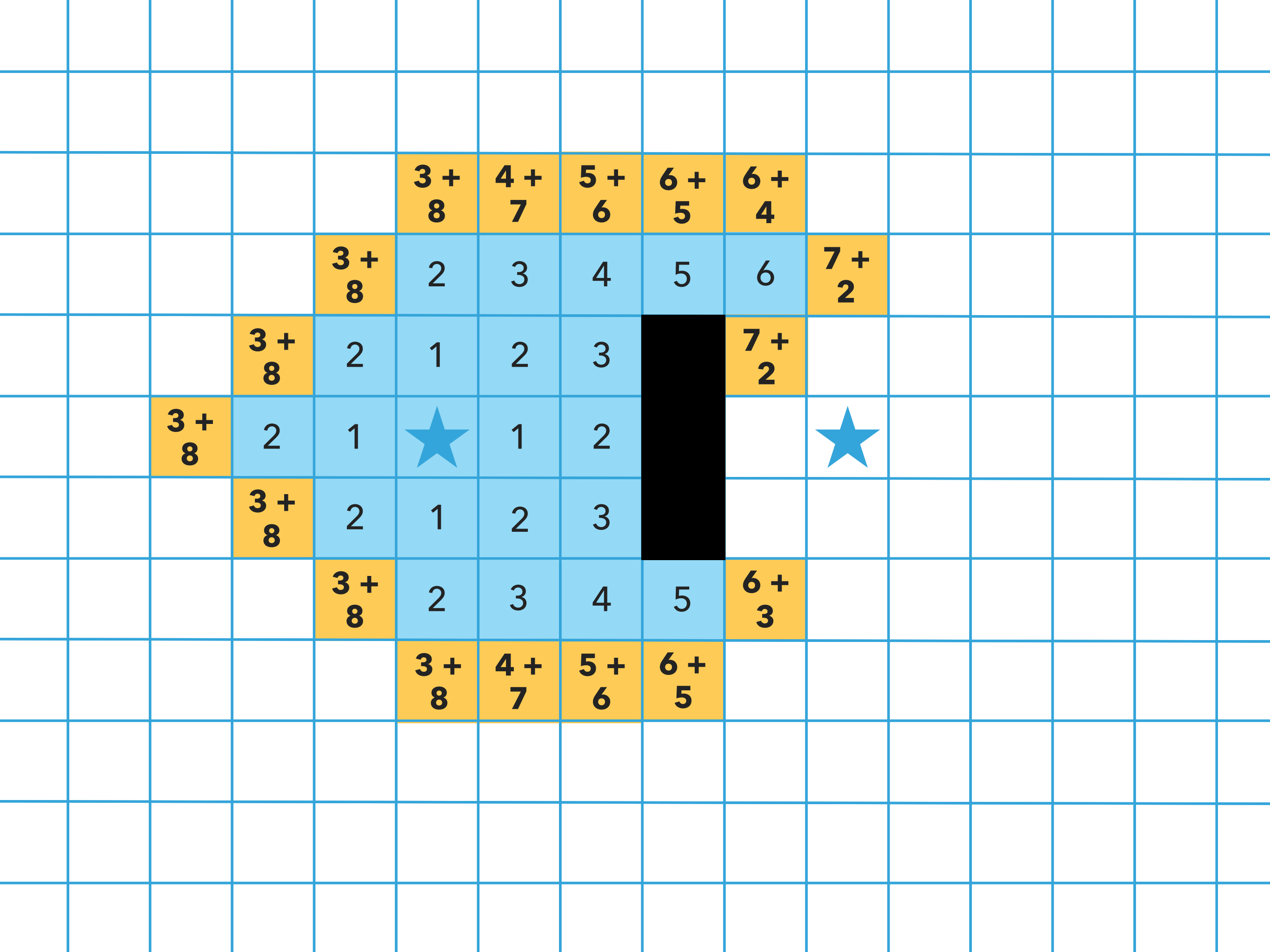


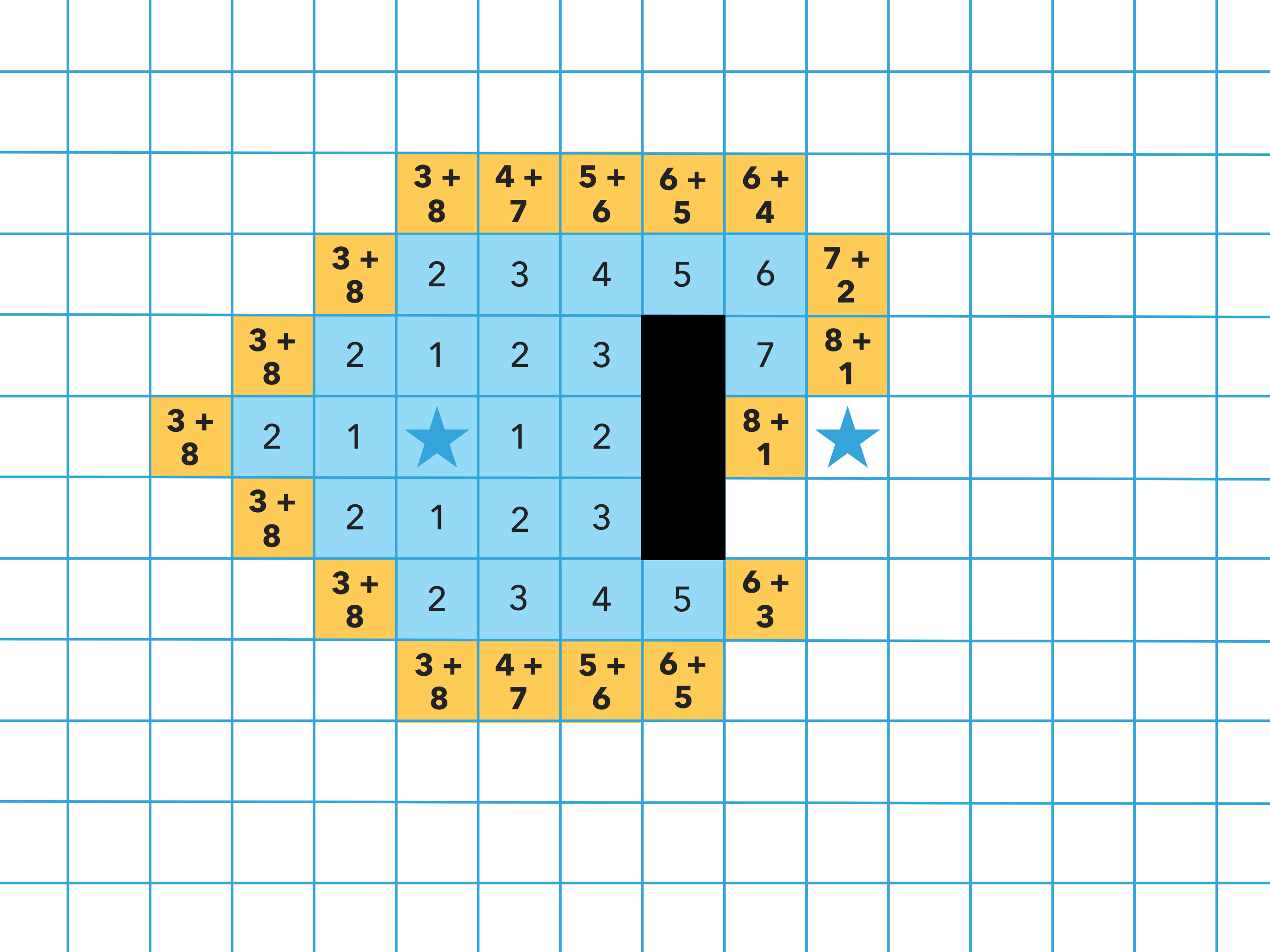


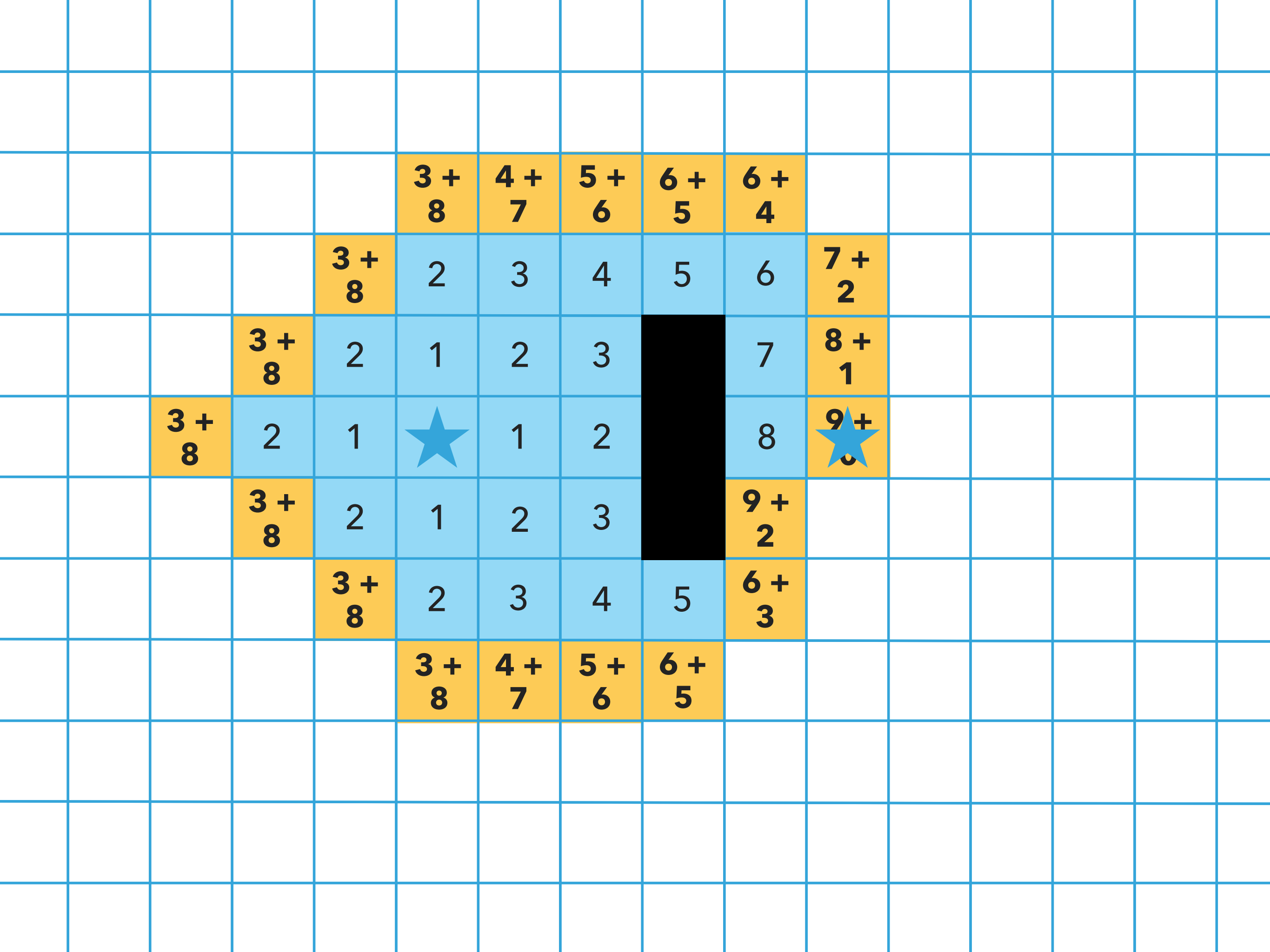


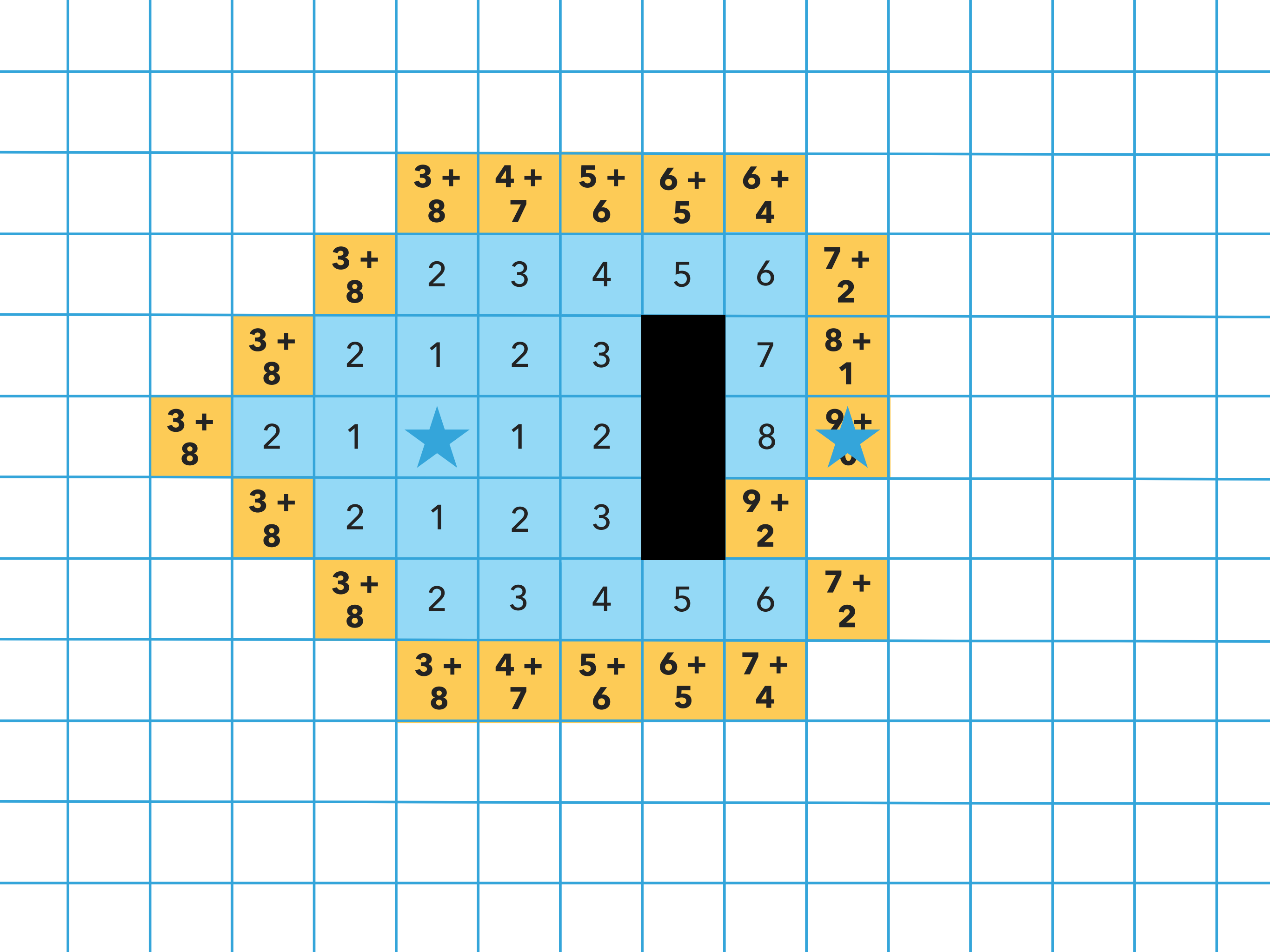


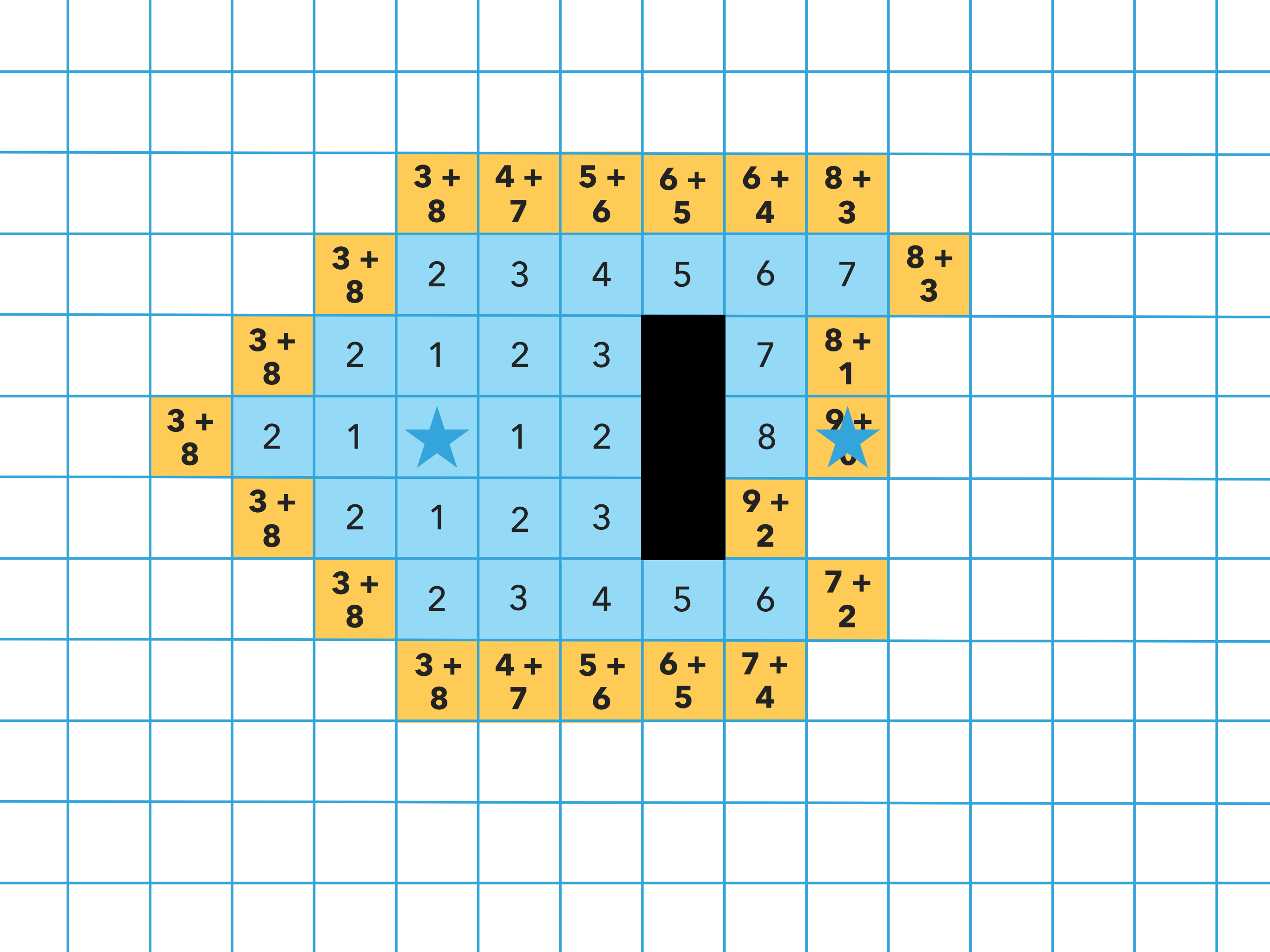


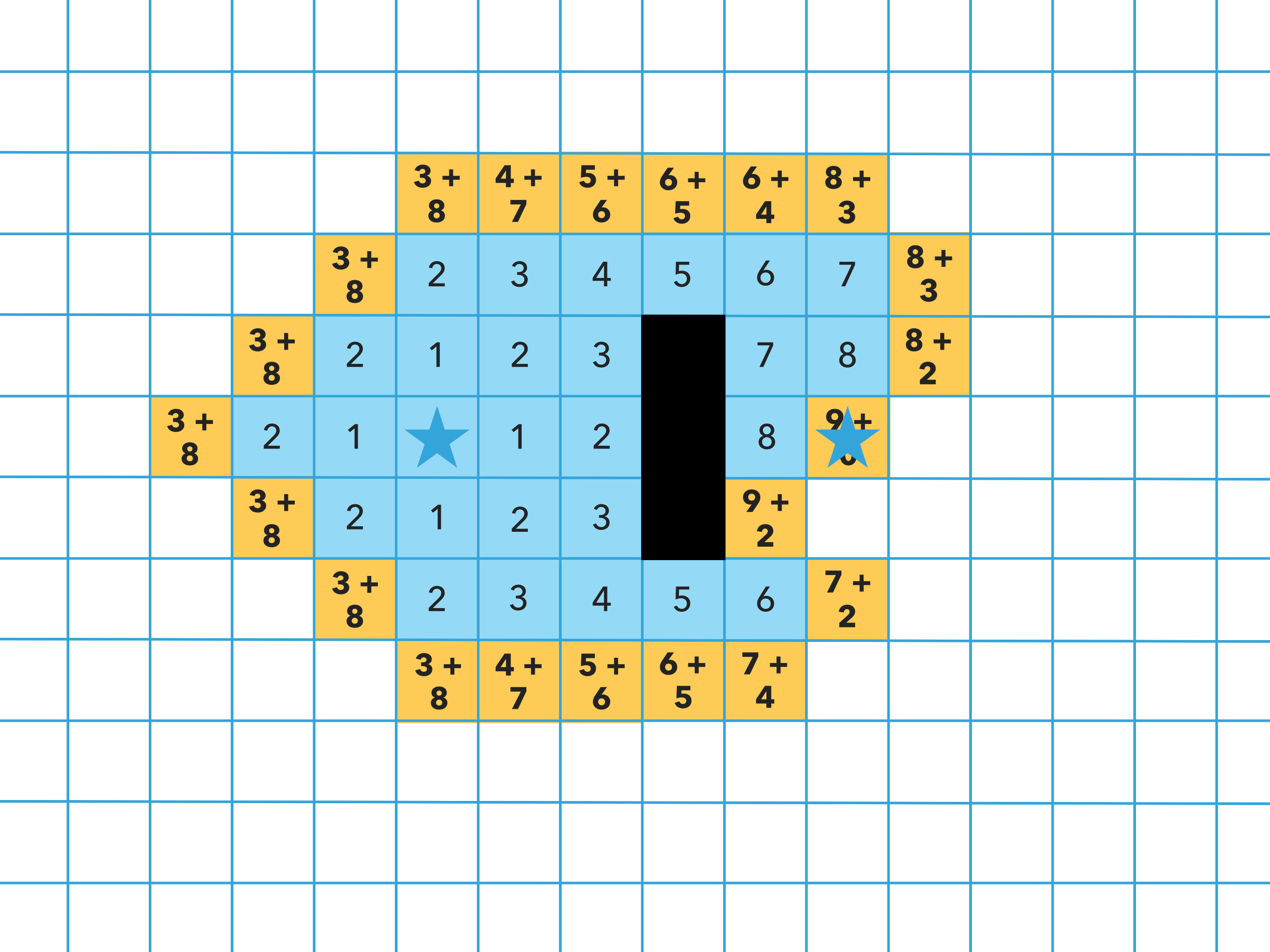


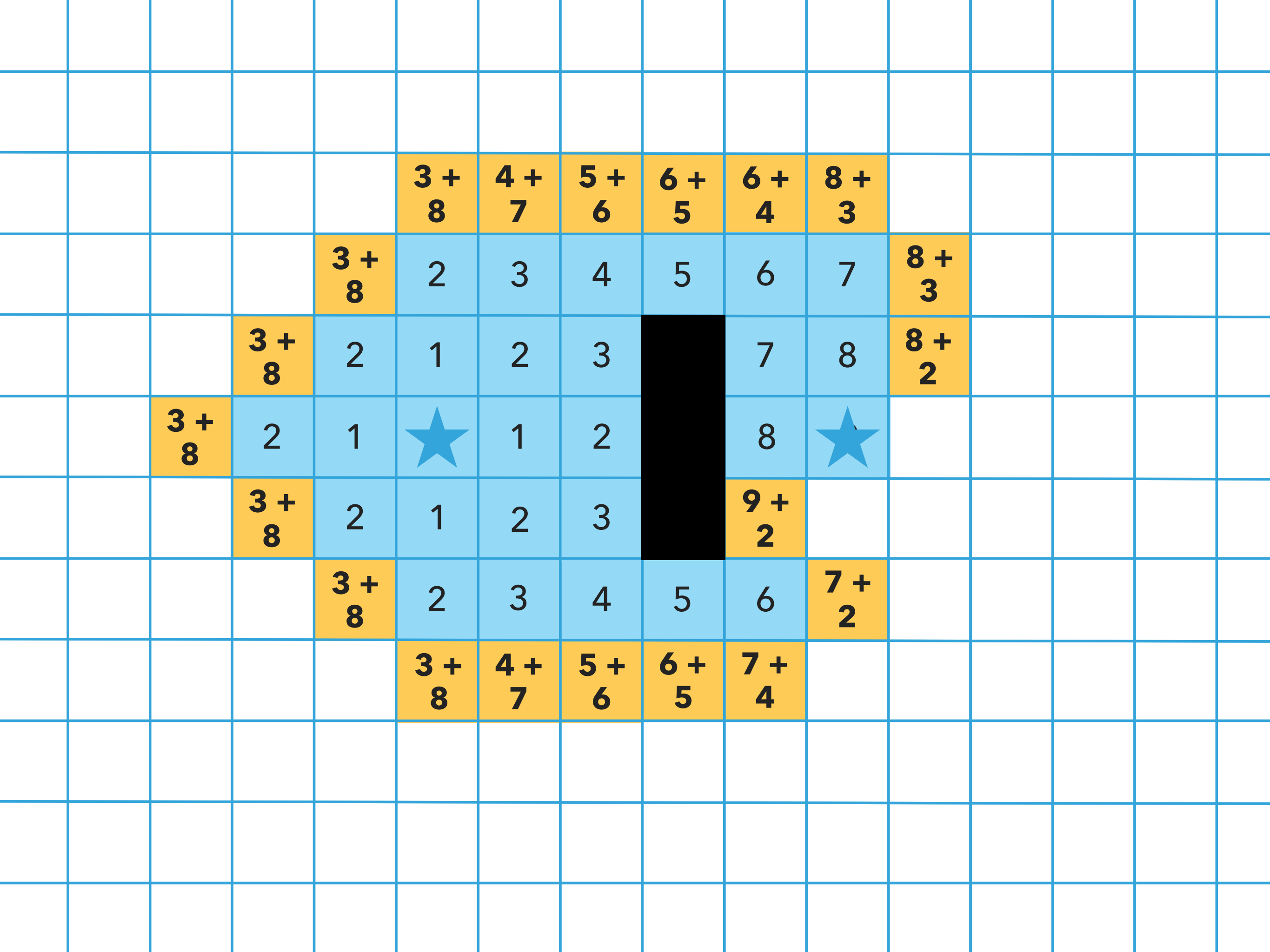


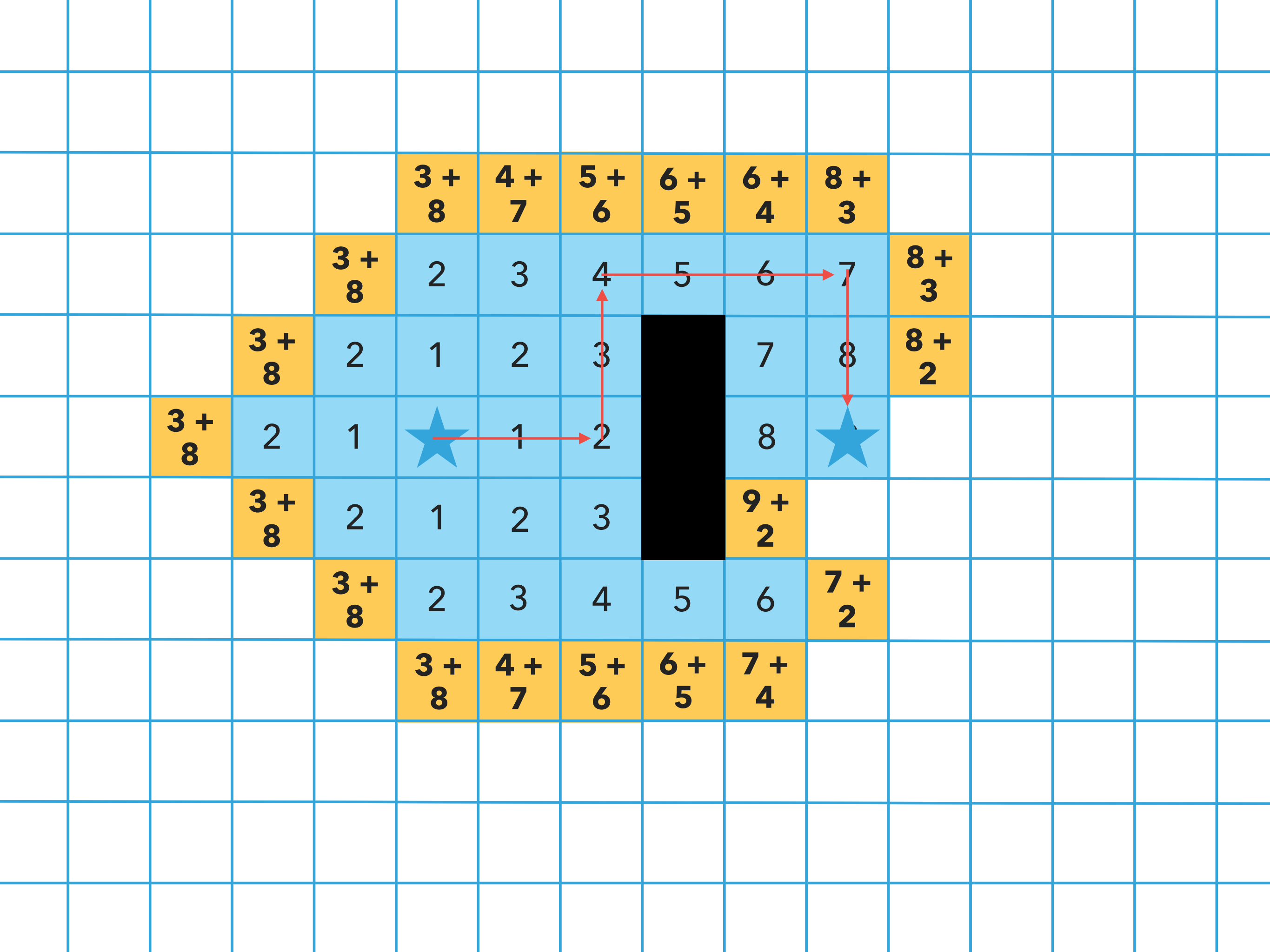


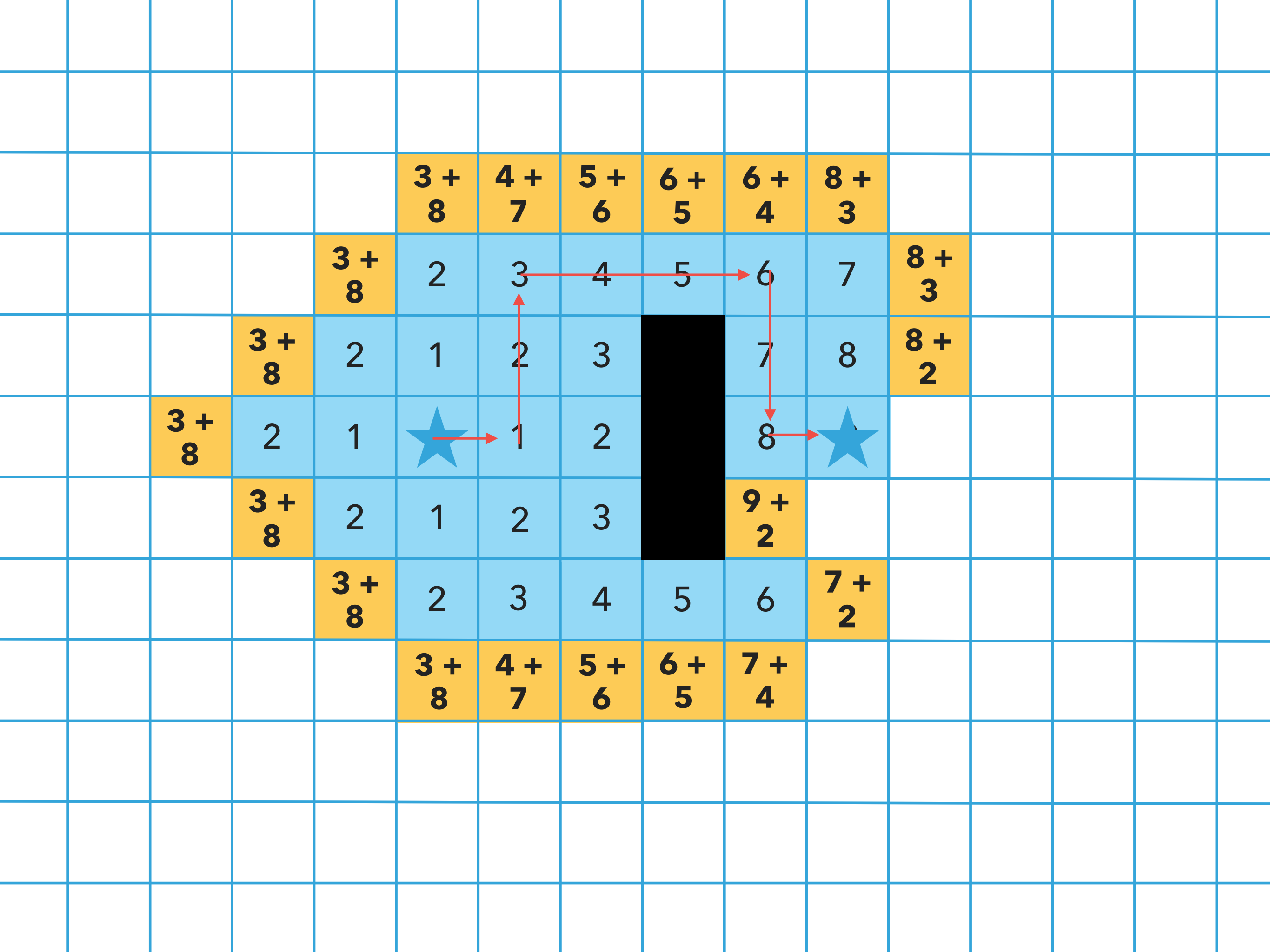


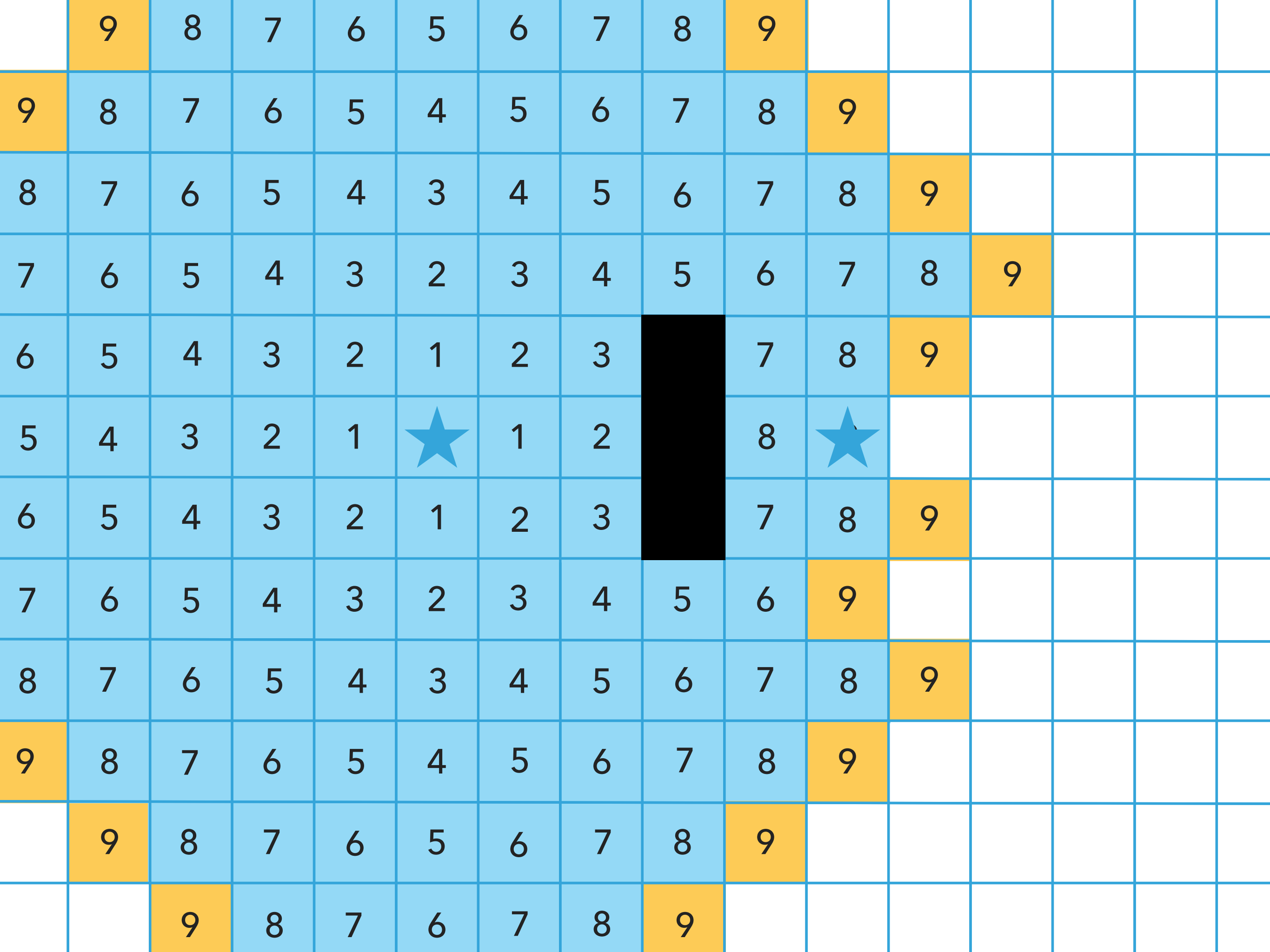












A* (PSEUDOCODE)

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- ▶ while q is not empty and end node isn't visited:
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 - ▶ $v = \text{last node of } p$
 - ▶ mark v as visited
 - ▶ for each unvisited neighbor:
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COMPARING DIJKSTRA AND A*

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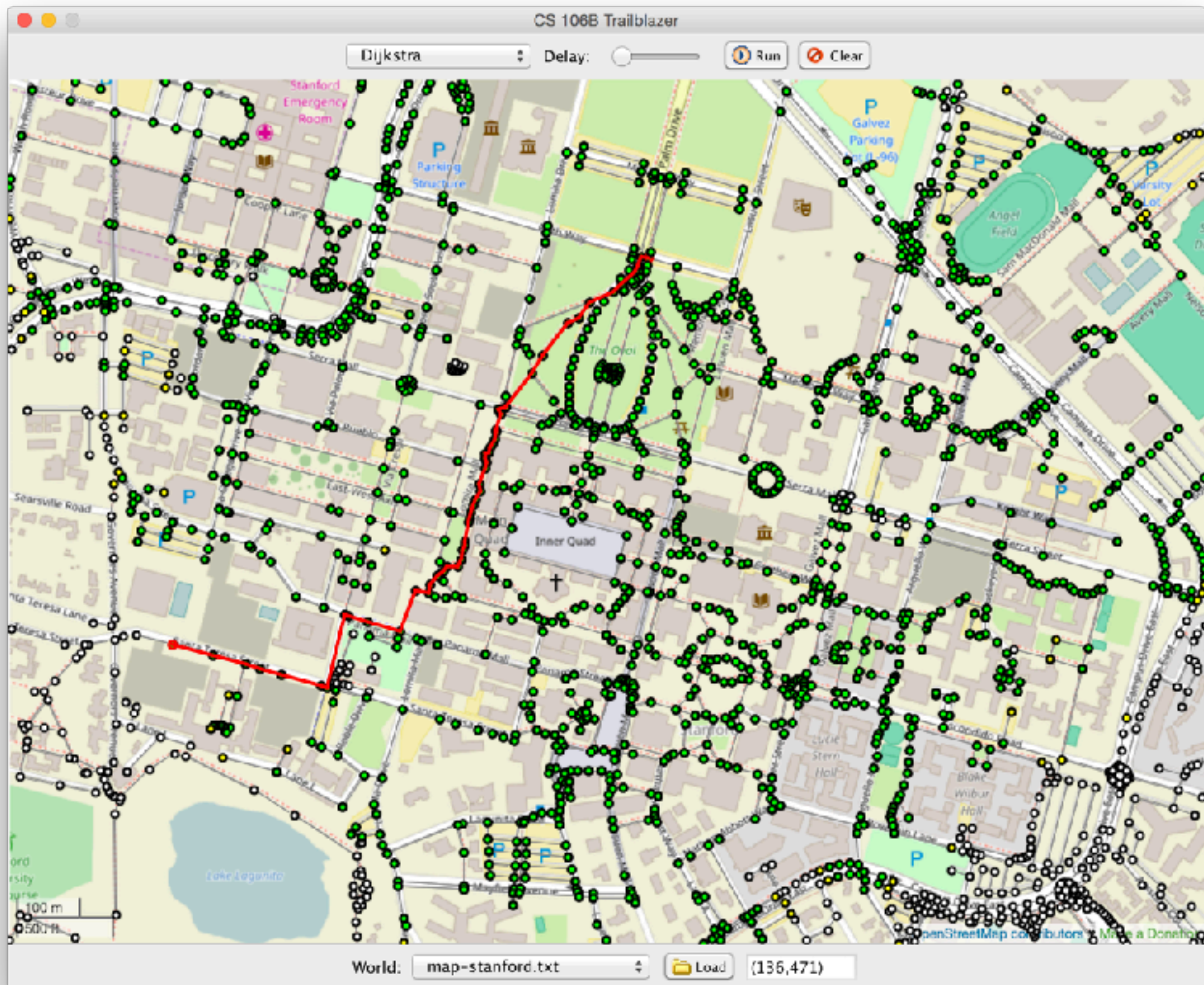
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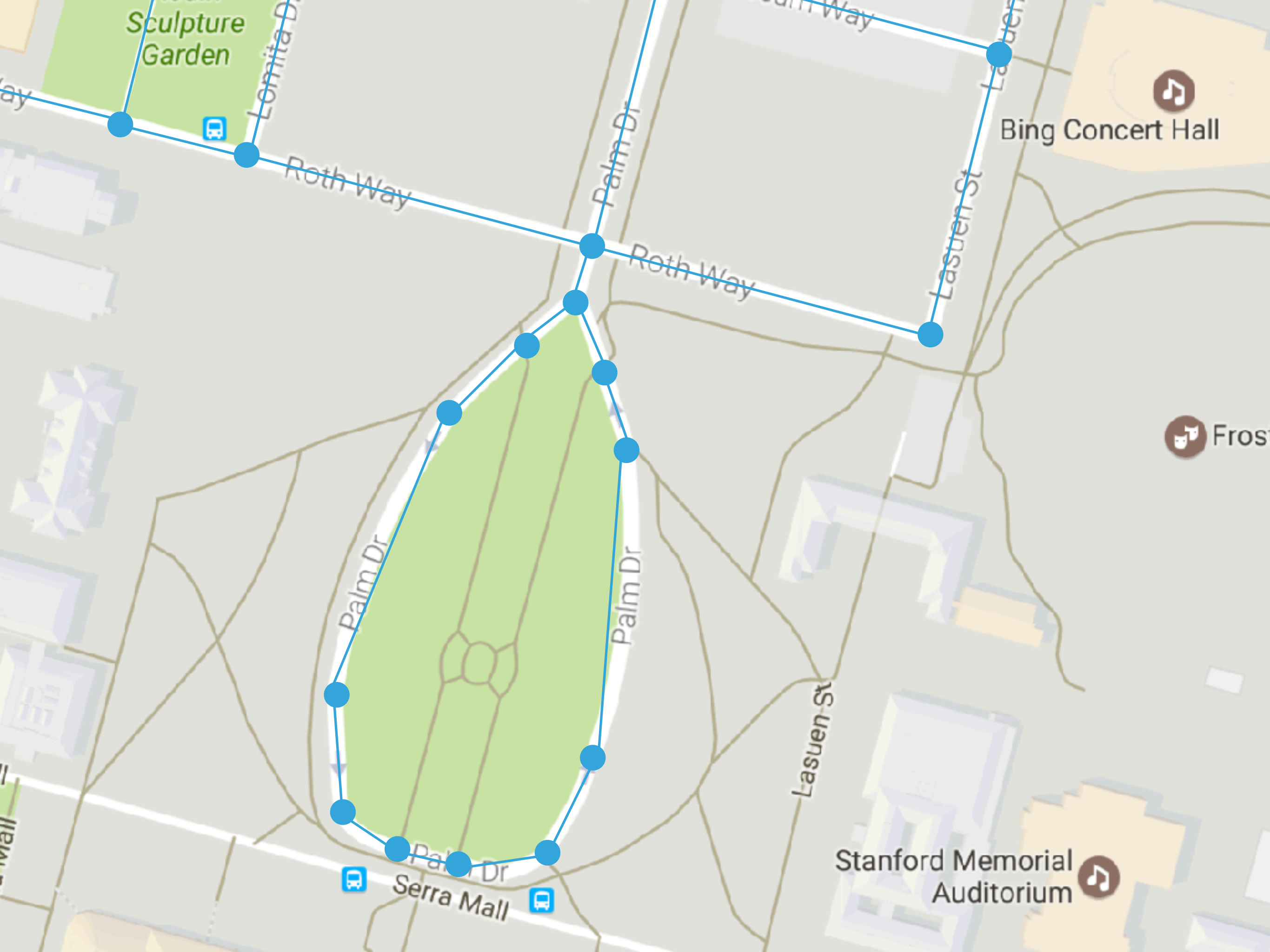
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**YOU WANT YOUR HEURISTIC TO BE AS LARGE
AS POSSIBLE**

**BUT YOU NEVER WANT IT TO BE LARGER THAN
THE ACTUAL COST.**

GOOGLE MAPS





Sculpture
Garden

Lomita Dr

Roth Way

Palm Dr

Roth Way

Bing Concert Hall

Lasuen St



Fros

Palm Dr

Palm Dr

Lasuen St

Serra Mall

Stanford Memorial
Auditorium



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 - ▶ Or... 190 million years

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 - ▶ Calculate the straight-line distance from A to B, and divide by the speed on the fastest highway

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- ▶ All of these and more?
 - ▶ You can use multiple heuristics and choose the max