

Programming Abstractions (Accelerated)

Winter 2017

Stanford University

Computer Science Department

THE LIFE CHANGING MAGIC OF

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DIJKSTRA AND A\*

Friday, March 10, 2017

Reading: Programming Abstractions in C++, Chapter 18.6

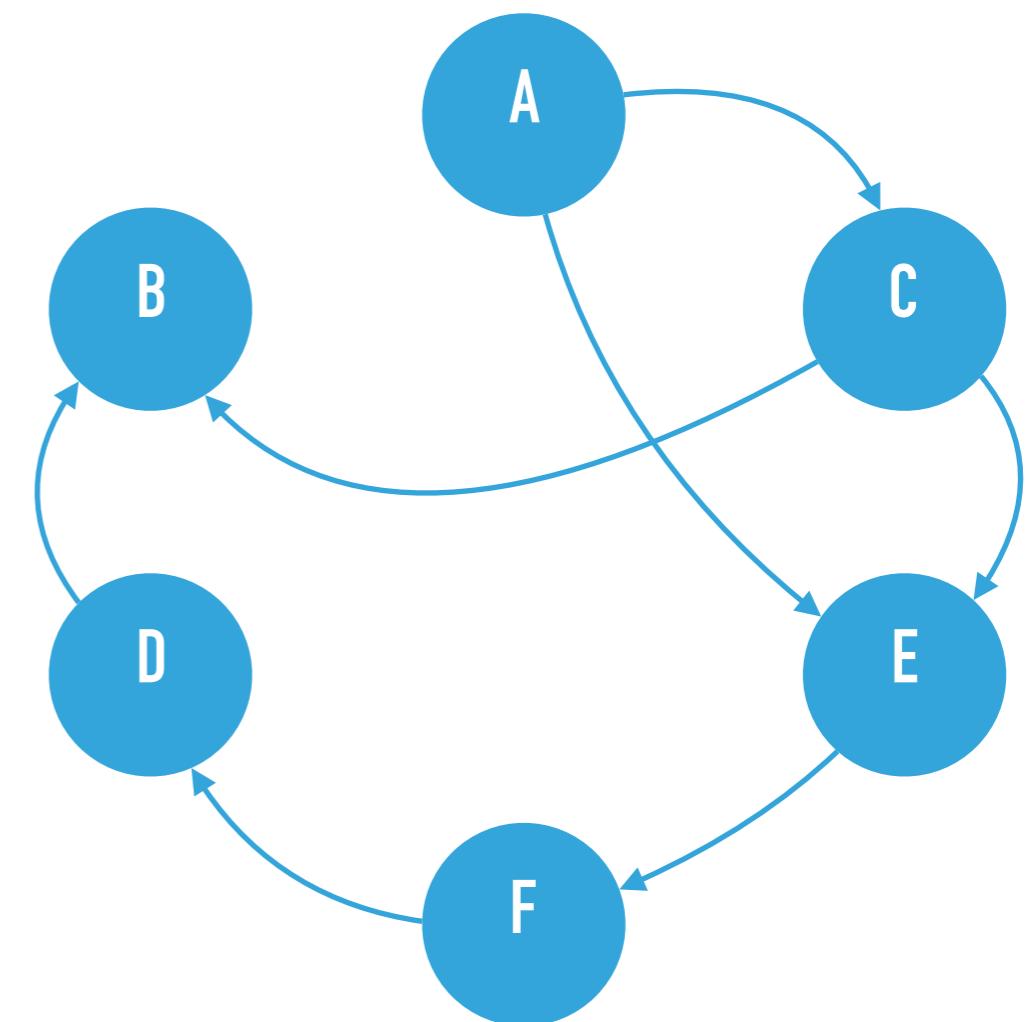
## TODAY'S TOPICS - MORE GRAPHS!

- ▶ Reviewing DFS and BFS
- ▶ Comparing DFS and BFS
- ▶ Making weighty decisions using Dijkstra's algorithm
- ▶ Looking into the future with A\*
- ▶ Google Maps

# REVIEWING DFS AND BFS

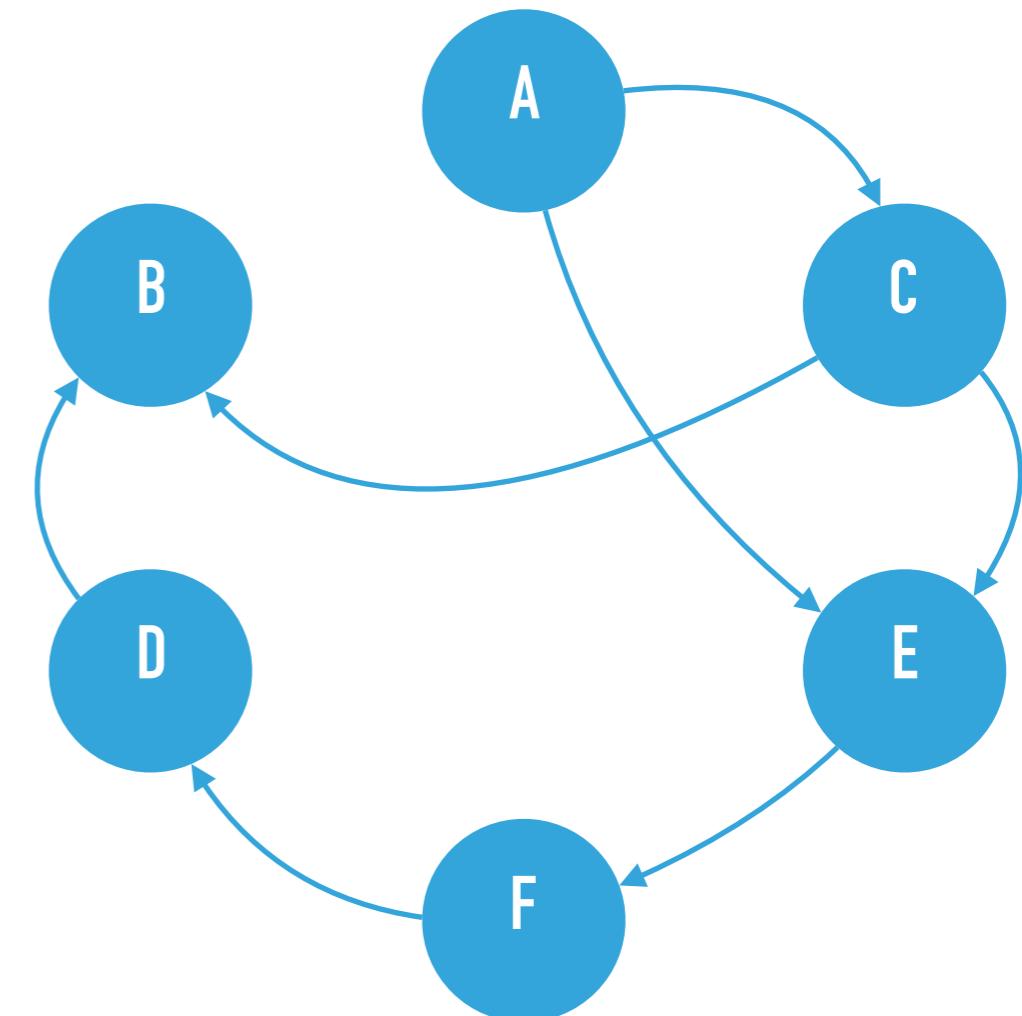
## DEPTH FIRST SEARCH

- ▶ Find a path from A to B using *iterative* depth first search
- ▶ (Assume that nodes are pushed onto the stack in *alphabetic order*)



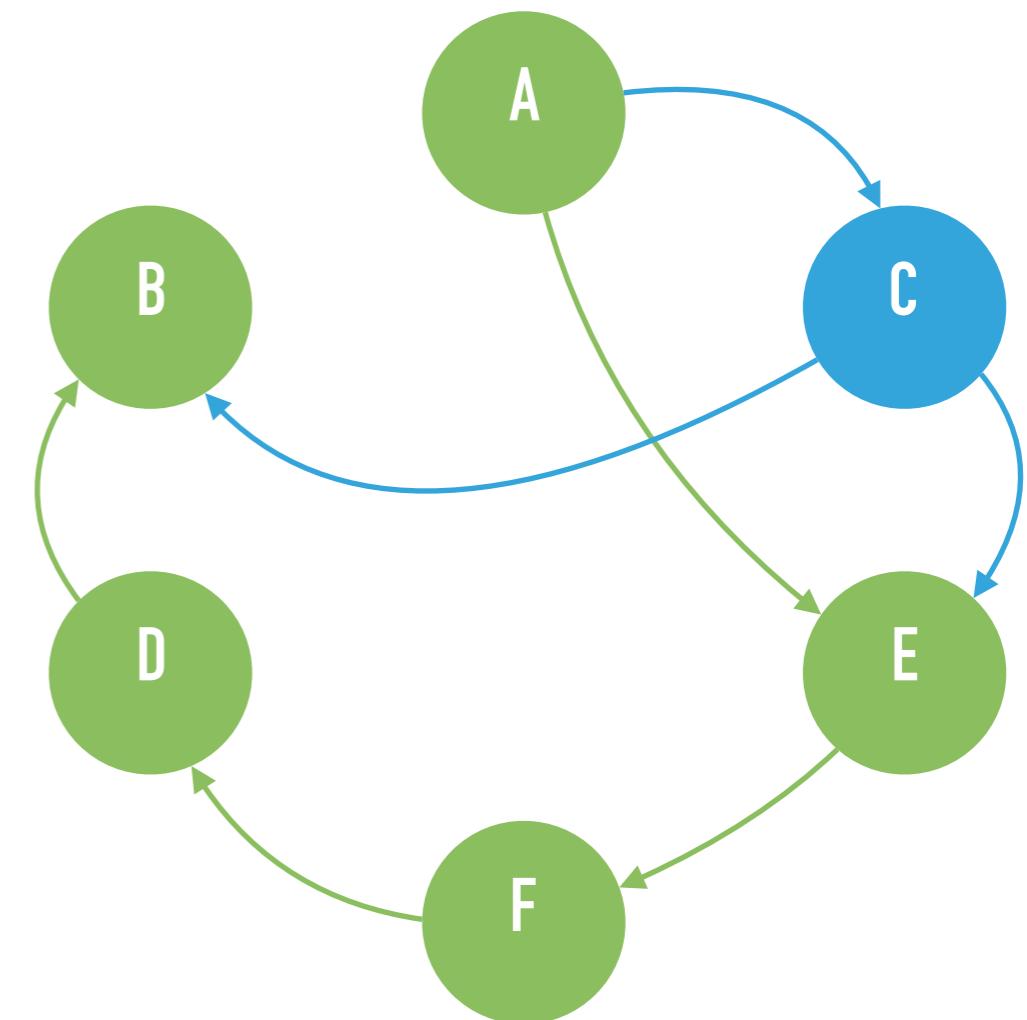
## DEPTH FIRST SEARCH (ITERATIVE PSEUDOCODE)

- ▶ create a path with just start node and push onto stack s
- ▶ while s is not empty
  - ▶  $p = s.pop()$
  - ▶  $v = \text{last node of } p$
  - ▶ if  $v$  is end, you're done
  - ▶ mark  $v$  as visited
  - ▶ for each unvisited neighbor:
    - ▶ create new path and append neighbor
    - ▶ push new path onto s



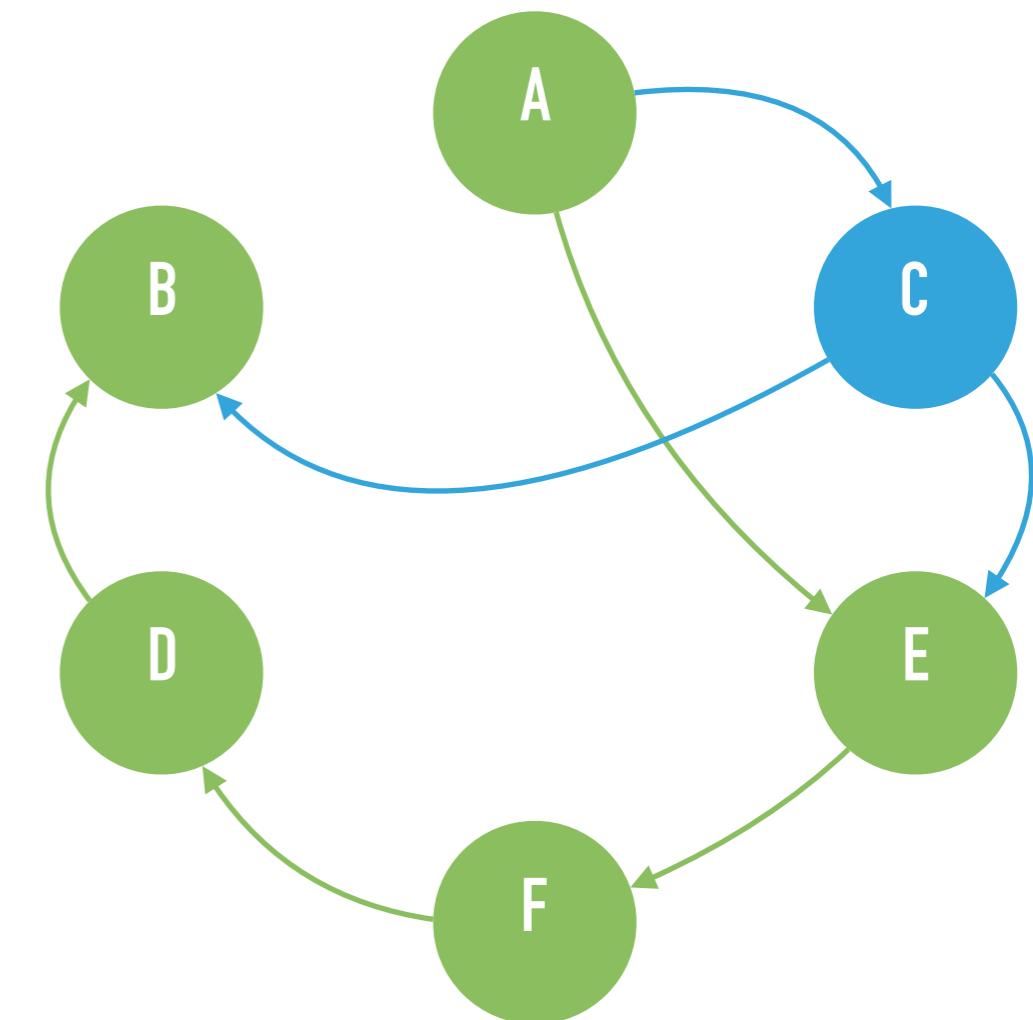
## DEPTH FIRST SEARCH

- ▶ Find a path from A to B using *iterative depth first search*
  - ▶ (Assume that nodes are pushed onto the stack in *alphabetic order*)
- ▶ A → E → F → D → B

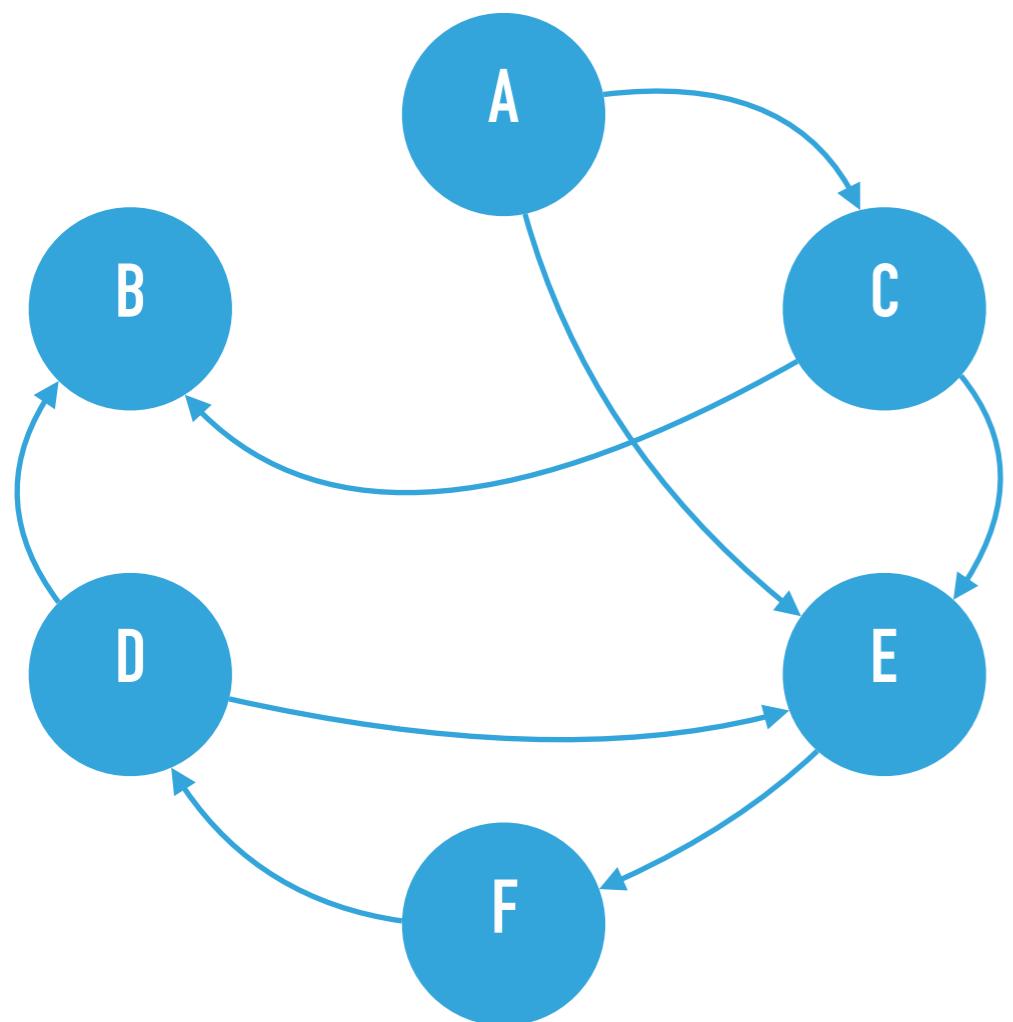


## DEPTH FIRST SEARCH

- ▶ Find a path from A to B using *iterative* depth first search
- ▶ (Assume that nodes are pushed onto the stack in *alphabetic order*)
- ▶ A → E → F → D → B
- ▶ Is this the shortest path?



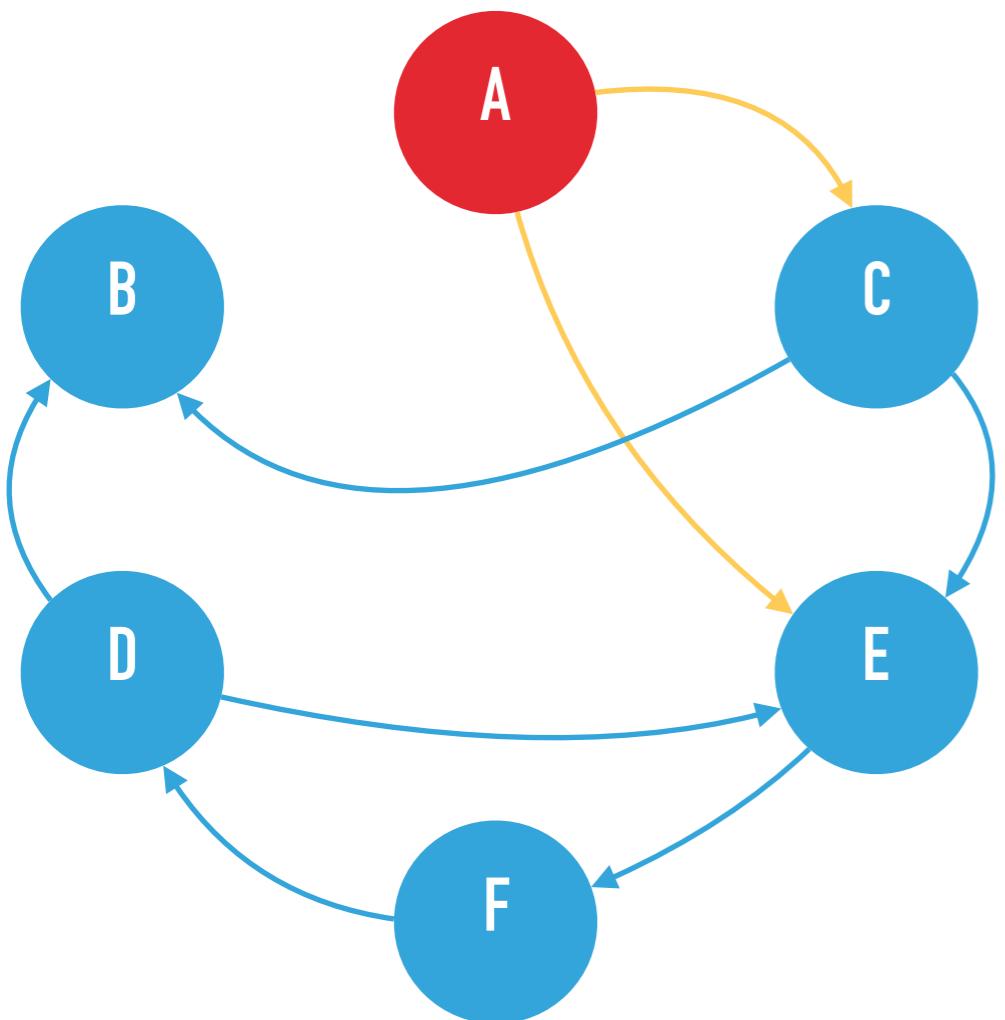
## DEPTH FIRST SEARCH



Paths to Consider (Stack)



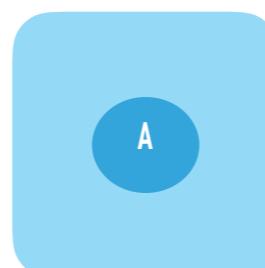
## DEPTH FIRST SEARCH



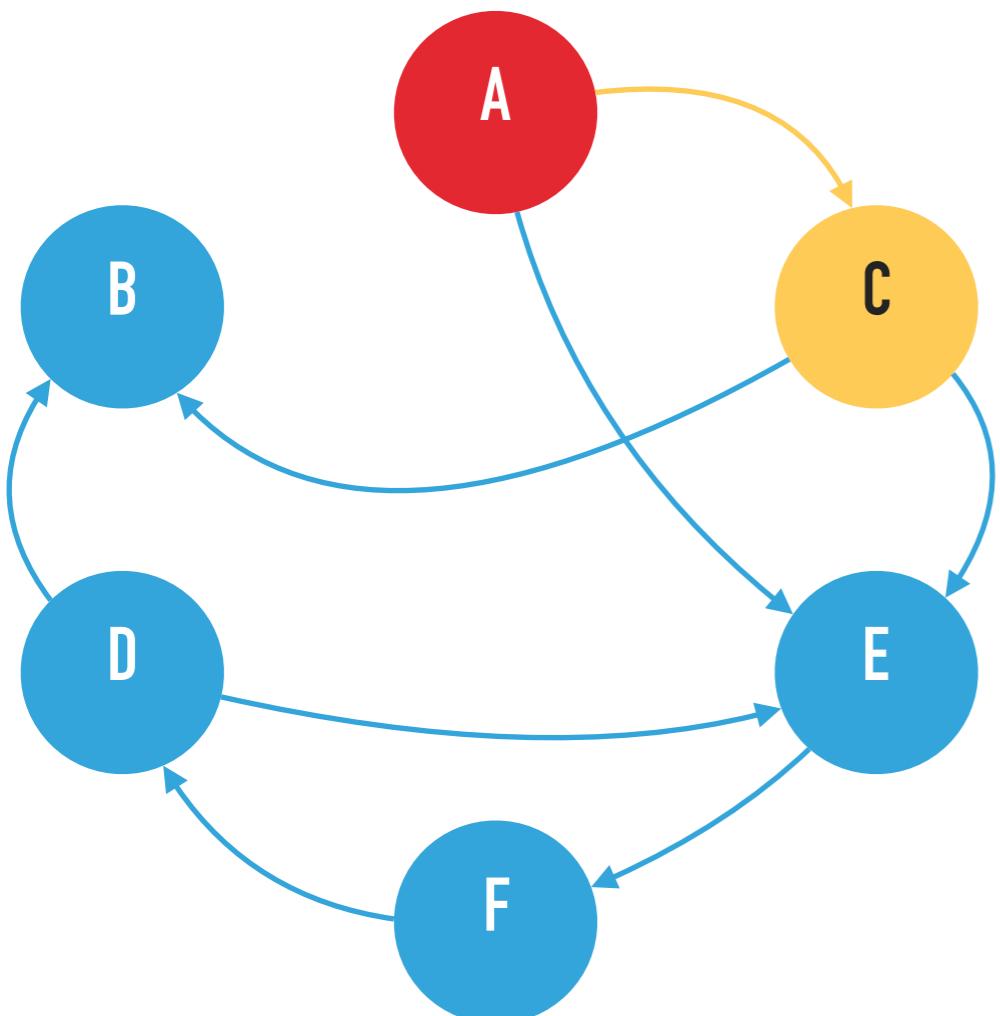
Paths to Consider (Stack)



Current Path



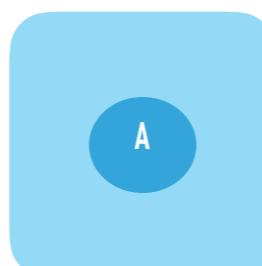
## DEPTH FIRST SEARCH



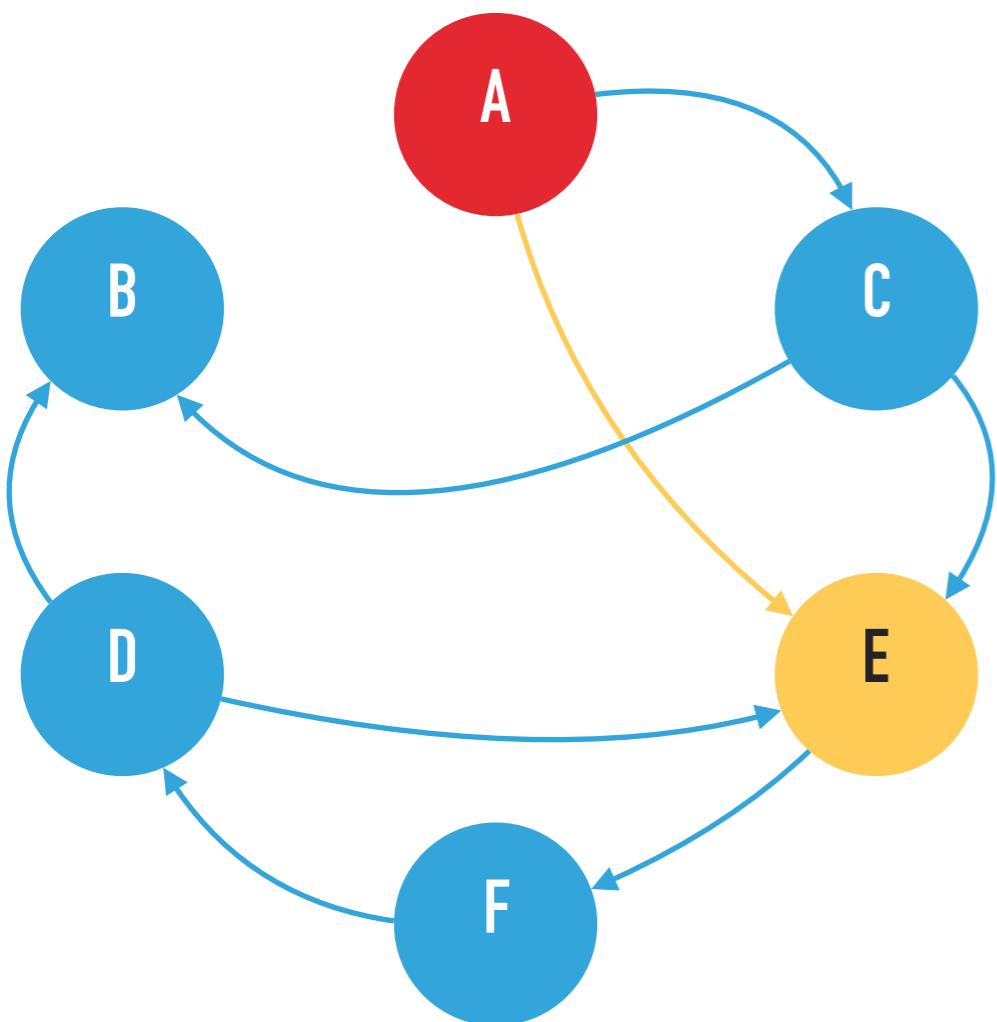
Paths to Consider (Stack)



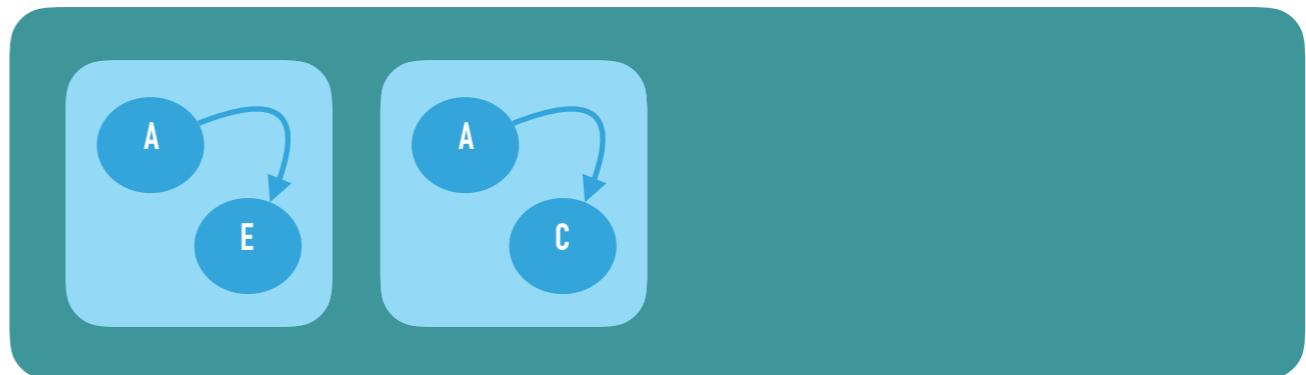
Current Path



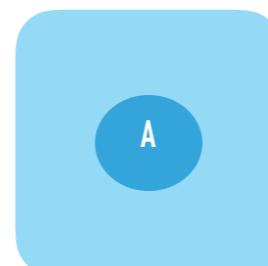
## DEPTH FIRST SEARCH



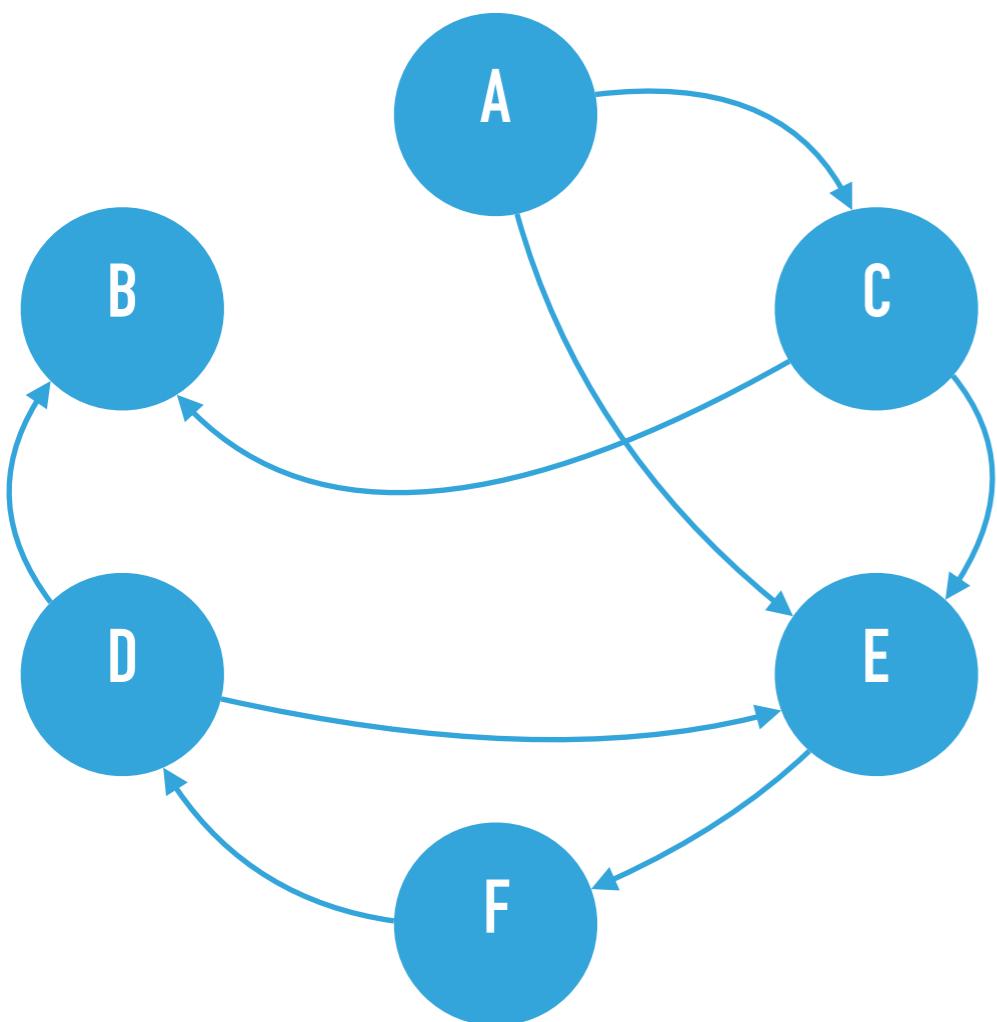
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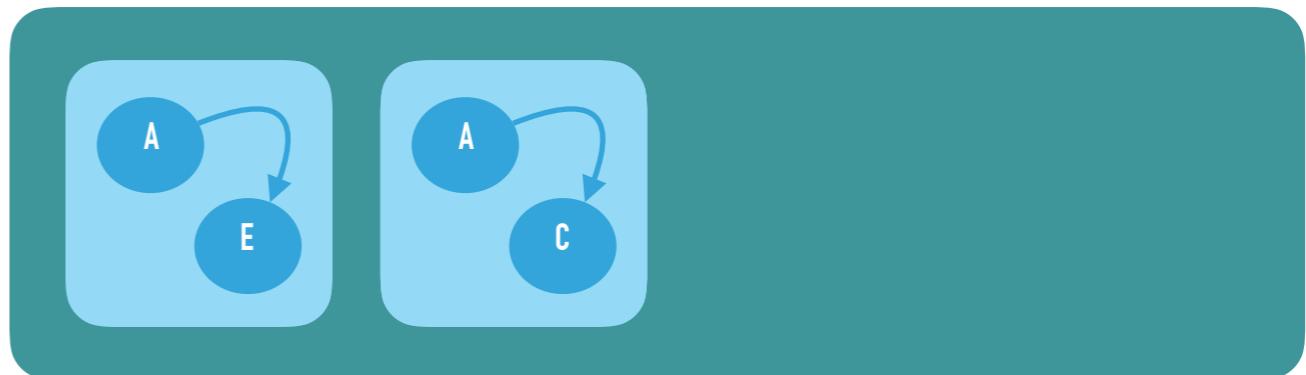
Current Path



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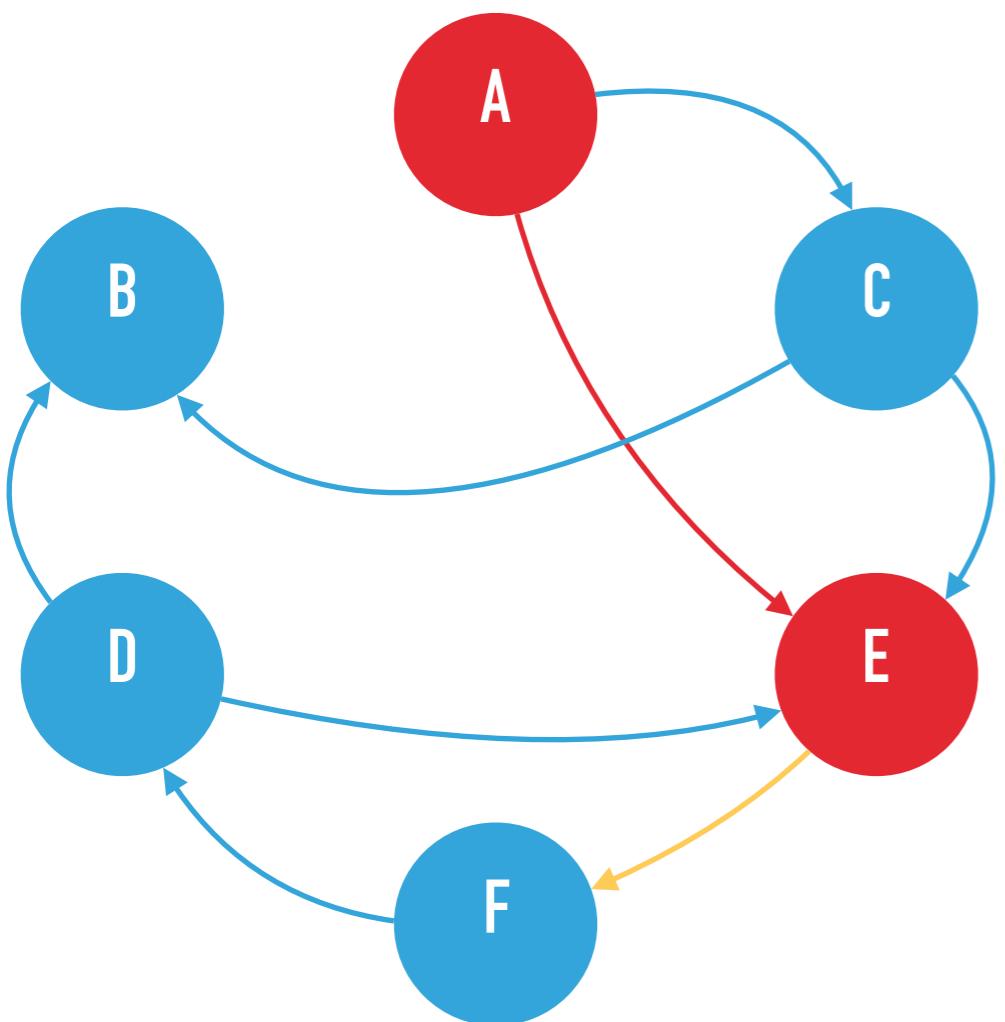


Paths to Consider (Stack)



Current Path

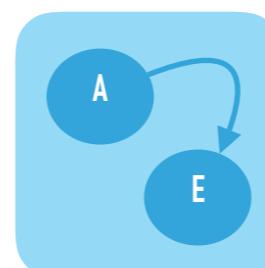
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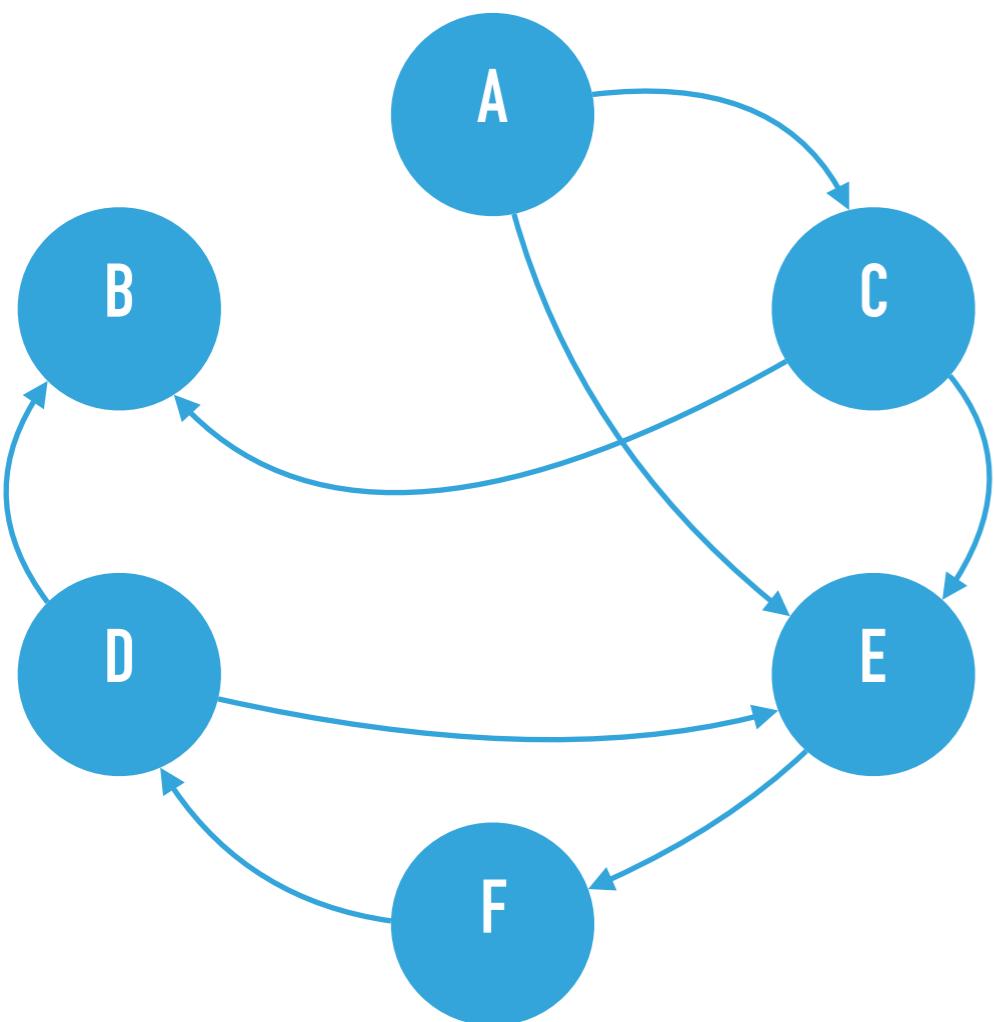
Paths to Consider (Stack)



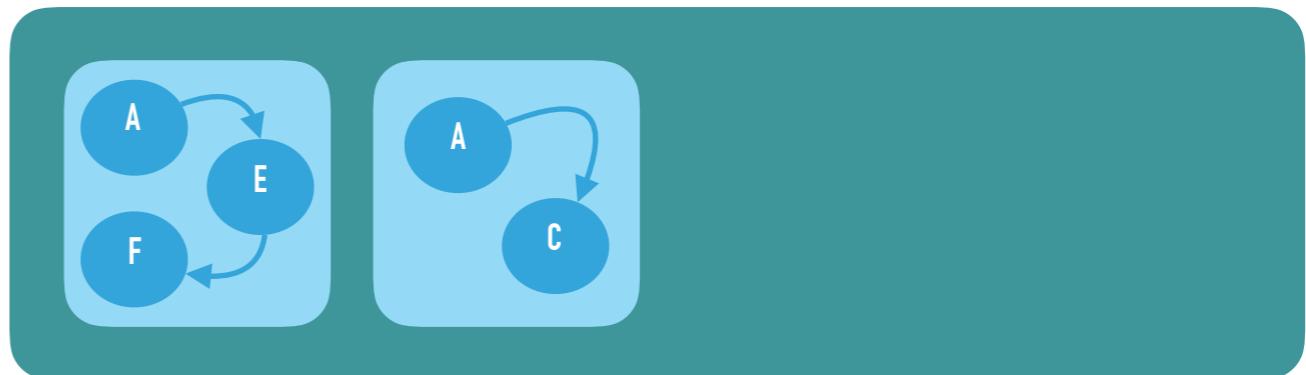
Current Path



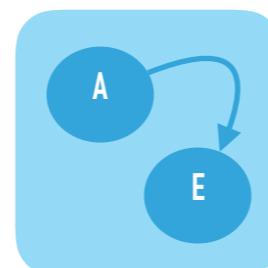
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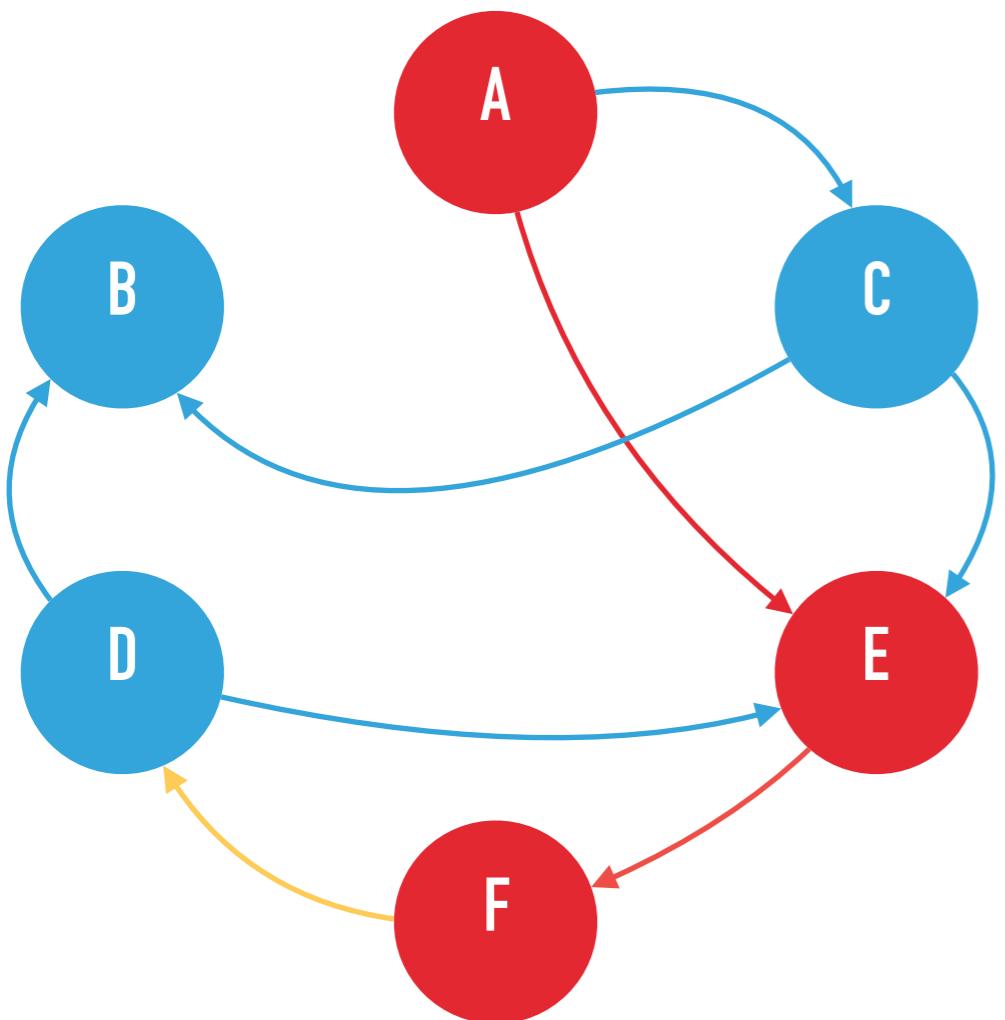
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Current Path



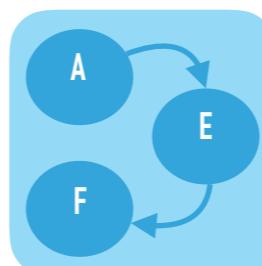
## DEPTH FIRST SEARCH



Paths to Consider (Stack)

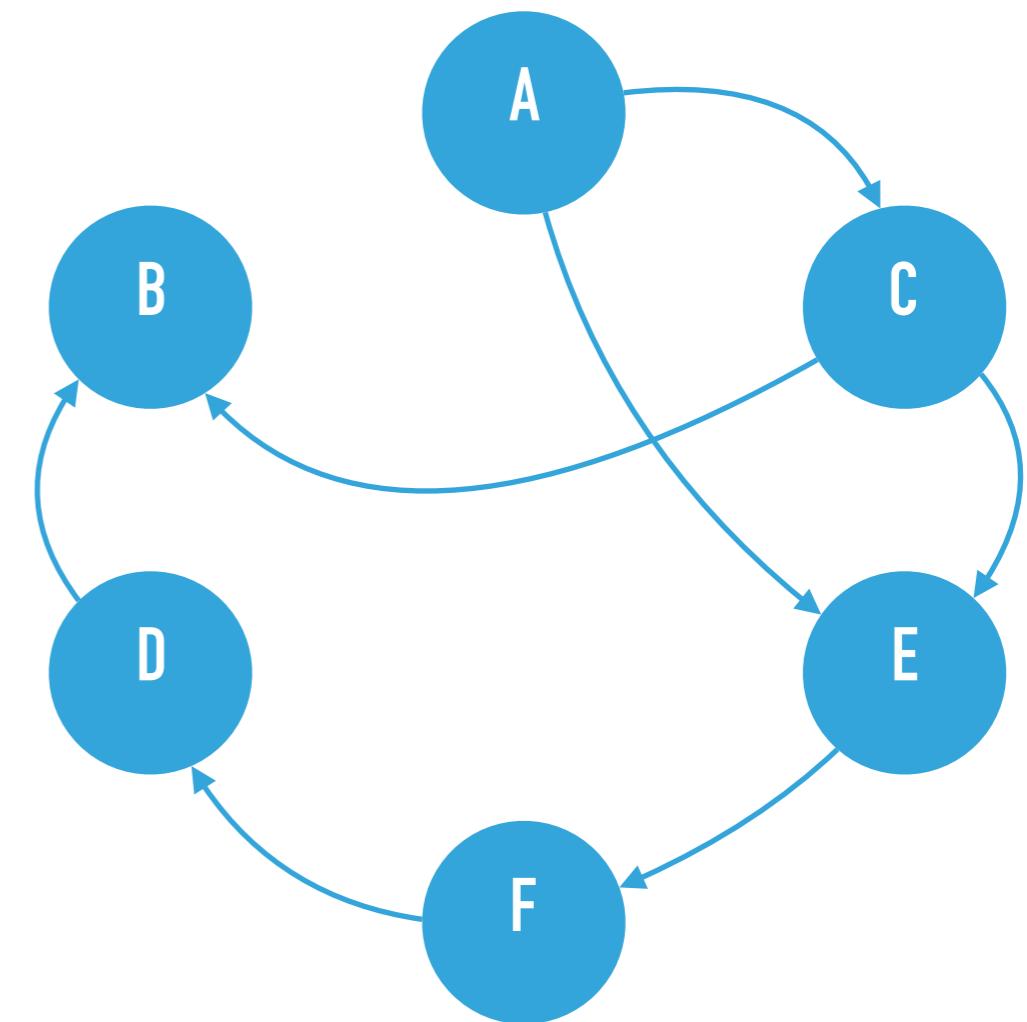


Current Path



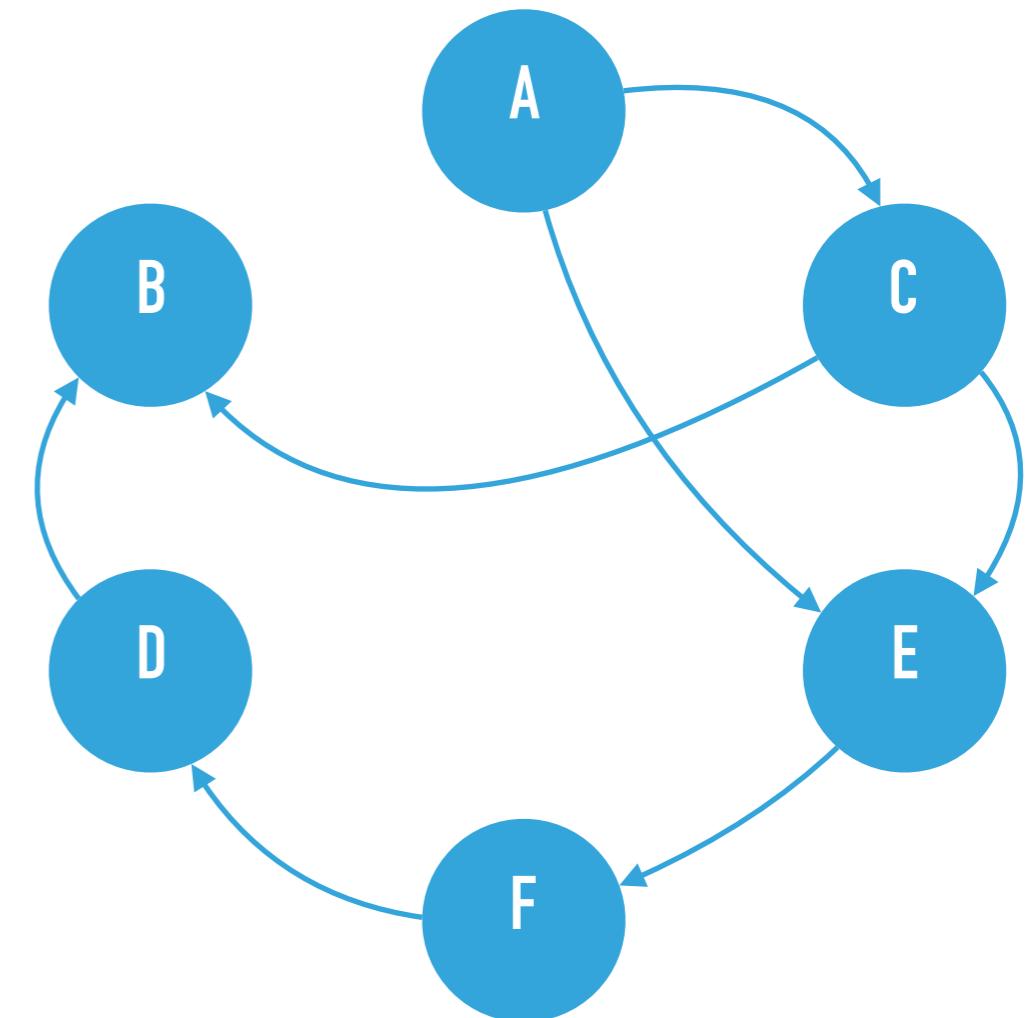
## BREADTH FIRST SEARCH

- ▶ Find a path from A to B using breadth first search
- ▶ (Assume that nodes are pushed onto the queue in *alphabetic order*)



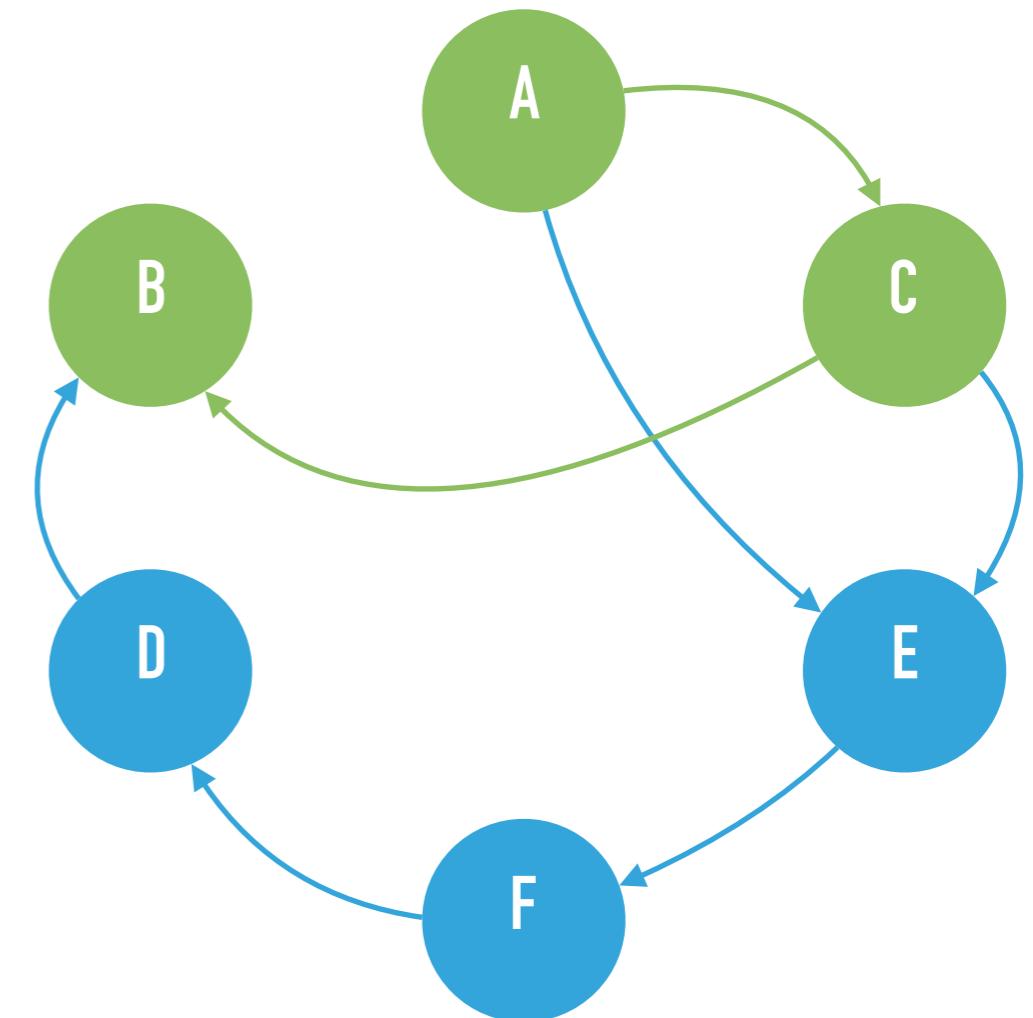
# BREADTH FIRST SEARCH (PSEUDOCODE)

- ▶ create a path with just start node and enqueue into queue q
- ▶ while q is not empty
  - ▶  $p = q.dequeue()$
  - ▶  $v = \text{last node of } p$
  - ▶ if  $v$  is end, you're done
  - ▶ mark  $v$  as visited
  - ▶ for each unvisited neighbor:
    - ▶ create new path and append neighbor
    - ▶ enqueue new path into q



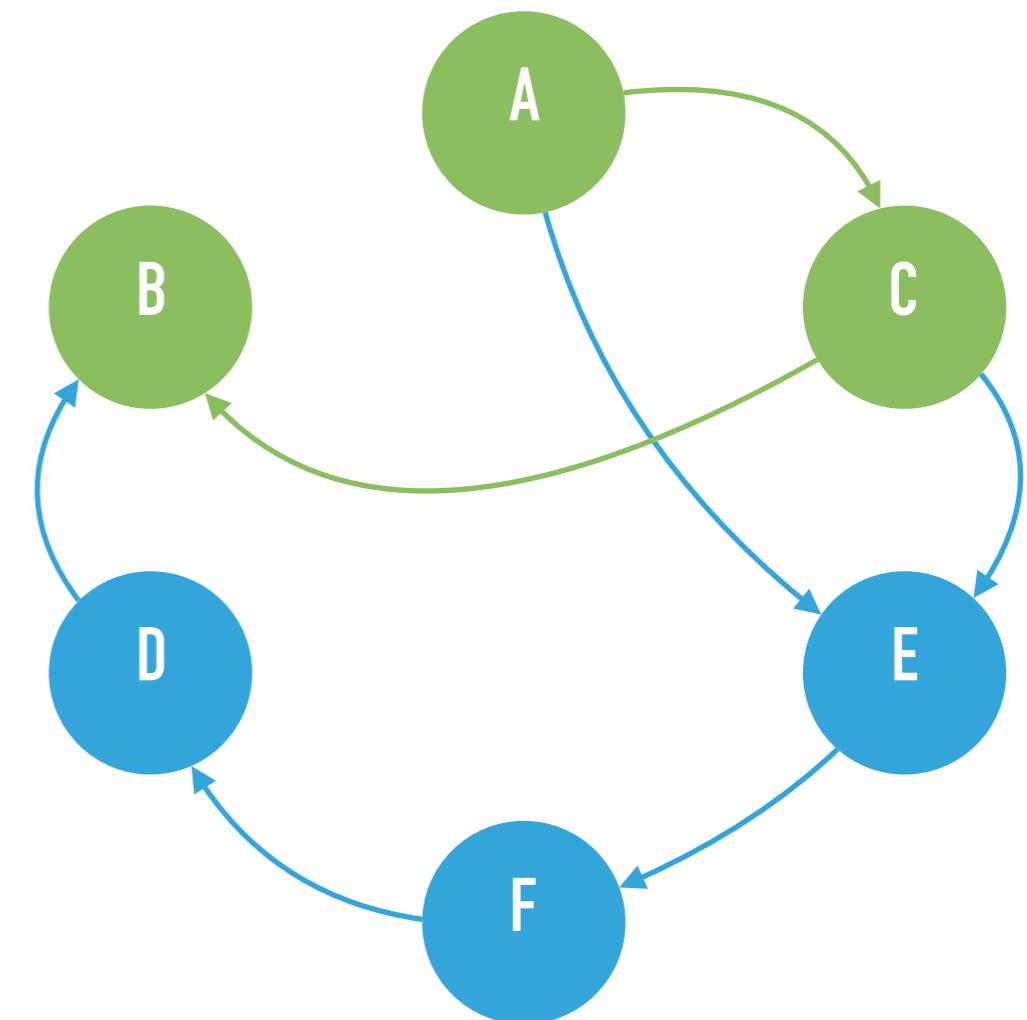
## BREADTH FIRST SEARCH

- ▶ Find a path from A to F using breadth first search
- ▶ (Assume that nodes are pushed onto the queue in *alphabetic order*)
- ▶ A → C → B



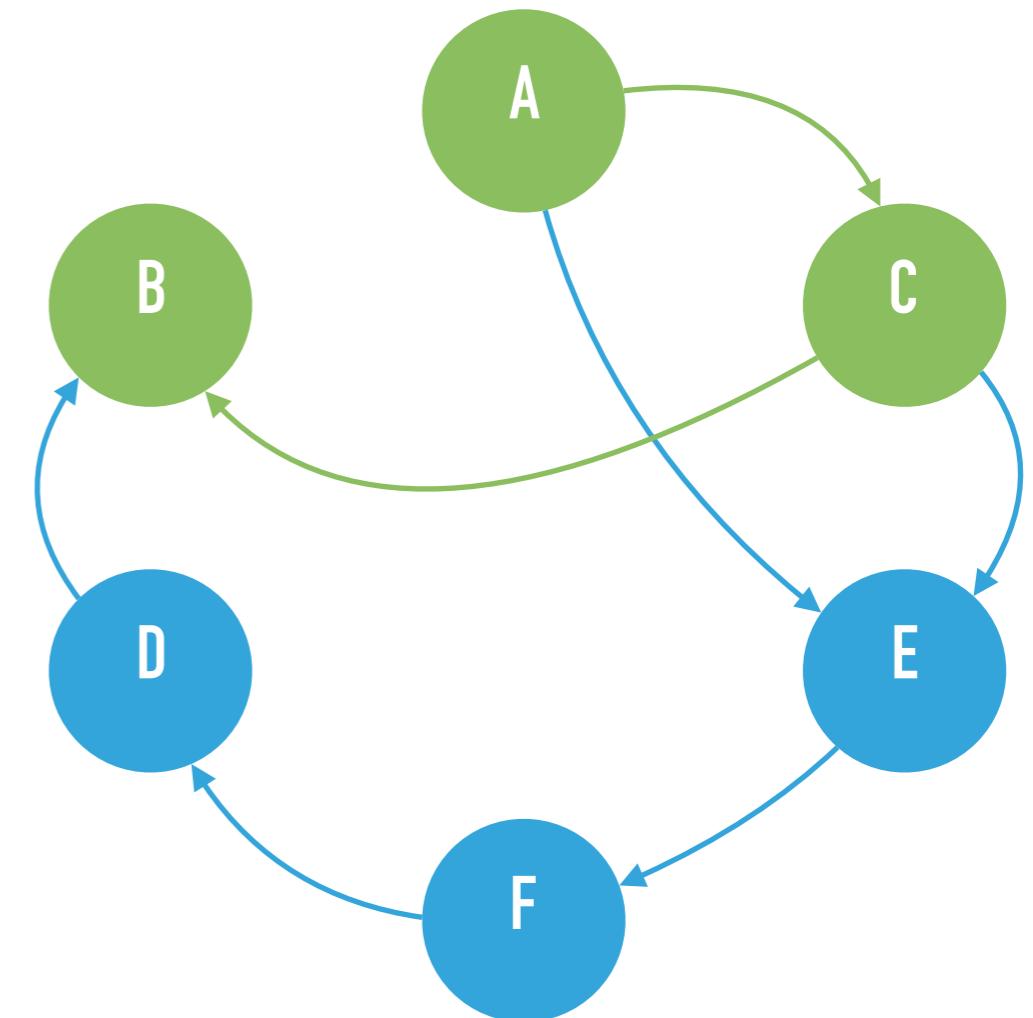
## BREADTH FIRST SEARCH

- ▶ Find a path from A to F using breadth first search
- ▶ (Assume that nodes are pushed onto the queue in *alphabetic order*)
- ▶ A → C → B
- ▶ Is *this* the shortest path?

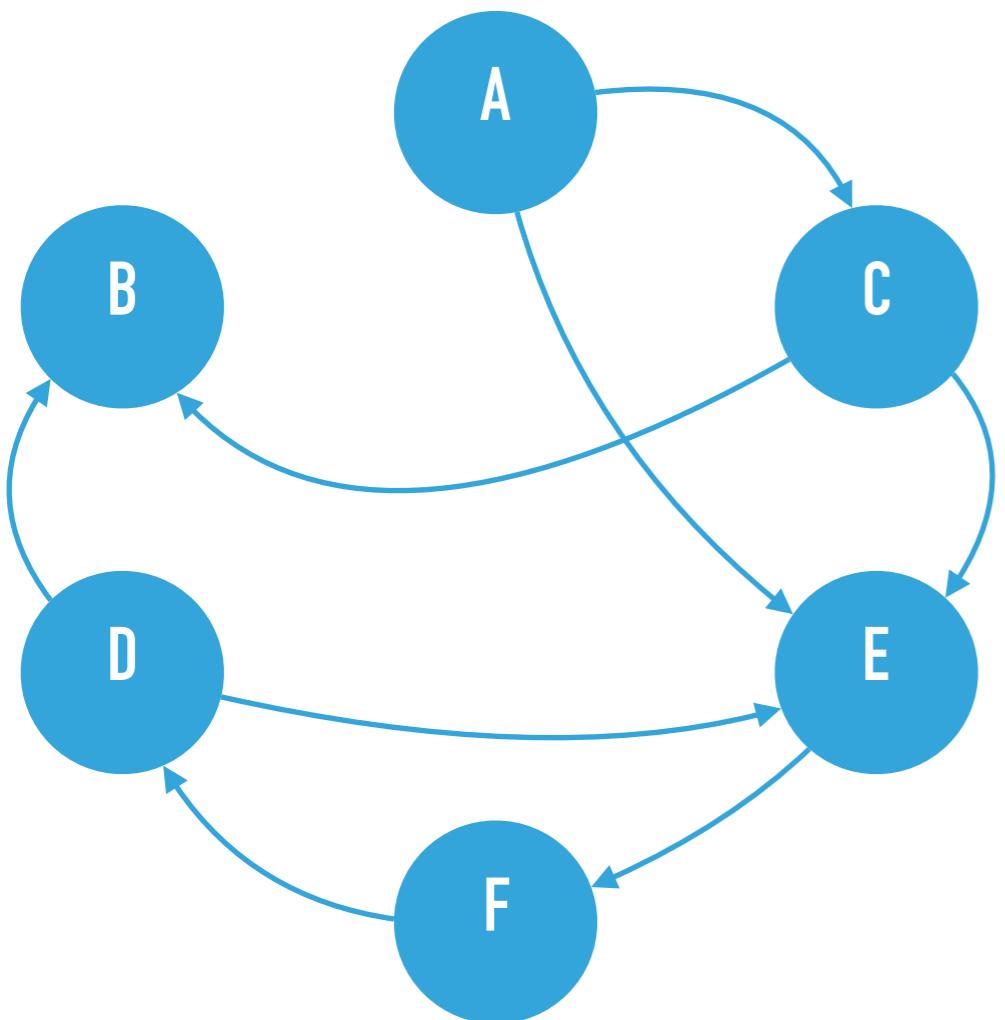


## BREADTH FIRST SEARCH

- ▶ Find a path from A to F using breadth first search
  - ▶ (Assume that nodes are pushed onto the queue in *alphabetic order*)
- ▶ A → C → B
- ▶ Is *this* the shortest path?
- ▶ Yes



## BREADTH FIRST SEARCH

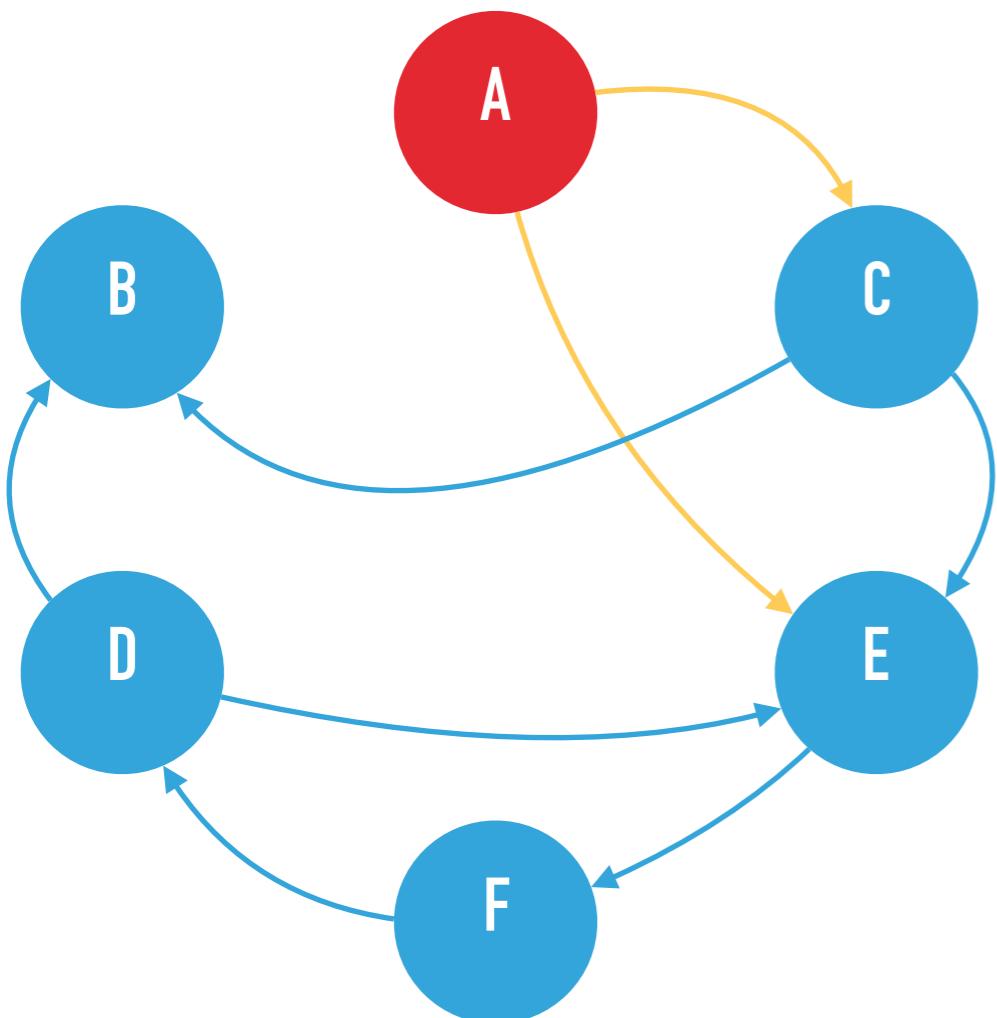


Paths to Consider (Queue)

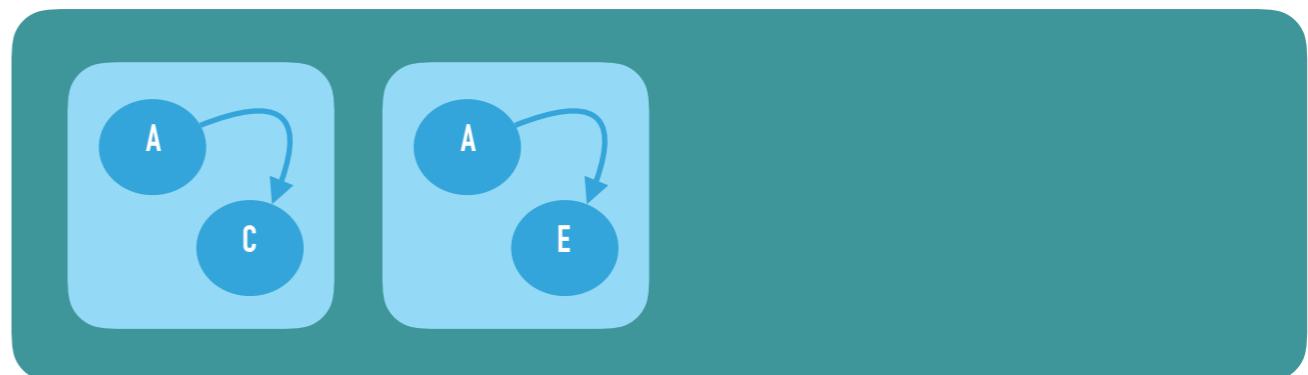


Current Path

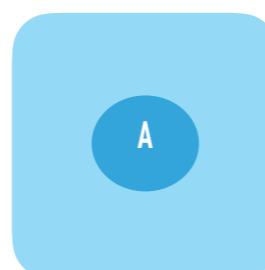
## BREADTH FIRST SEARCH



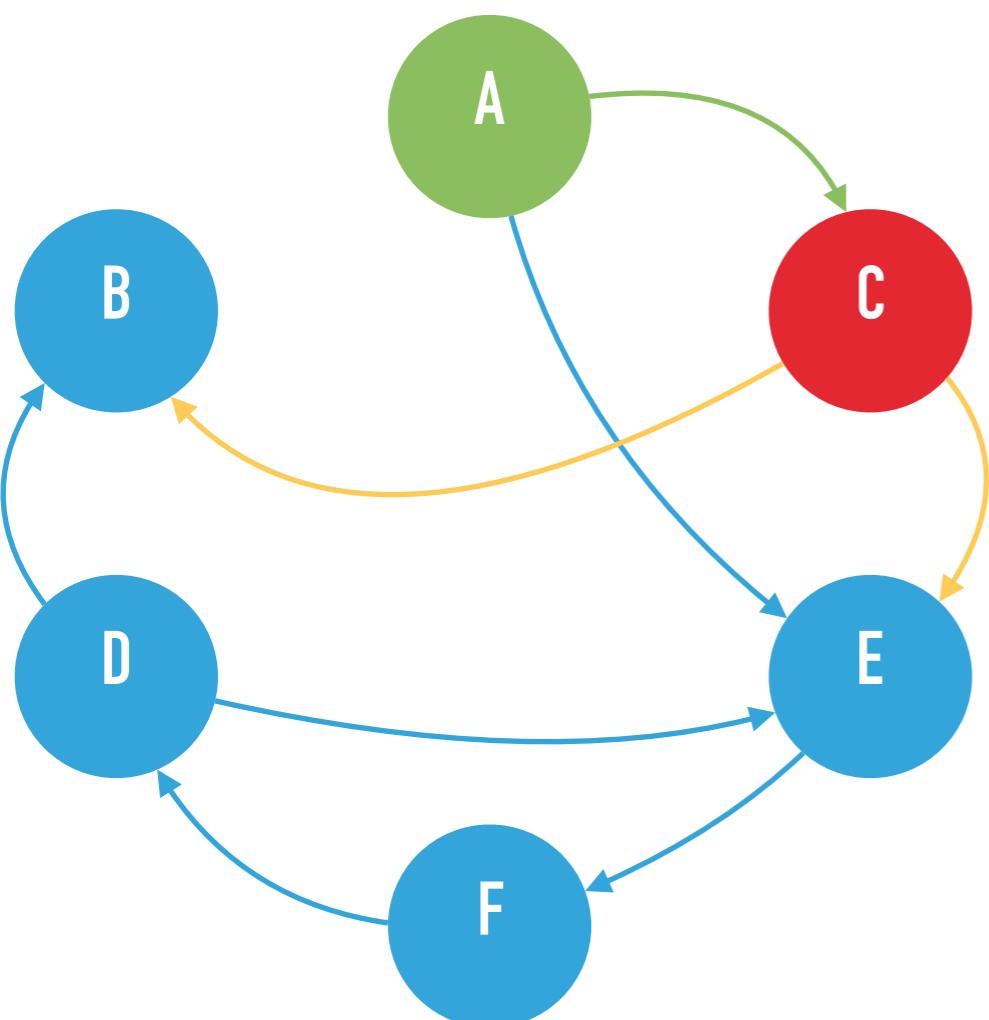
Paths to Consider (Queue)



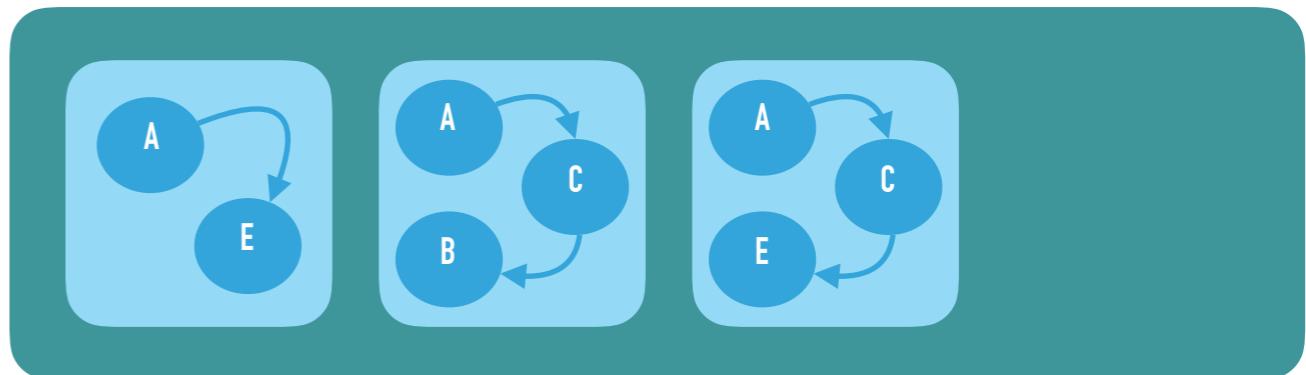
Current Path



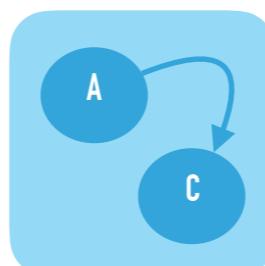
## BREADTH FIRST SEARCH



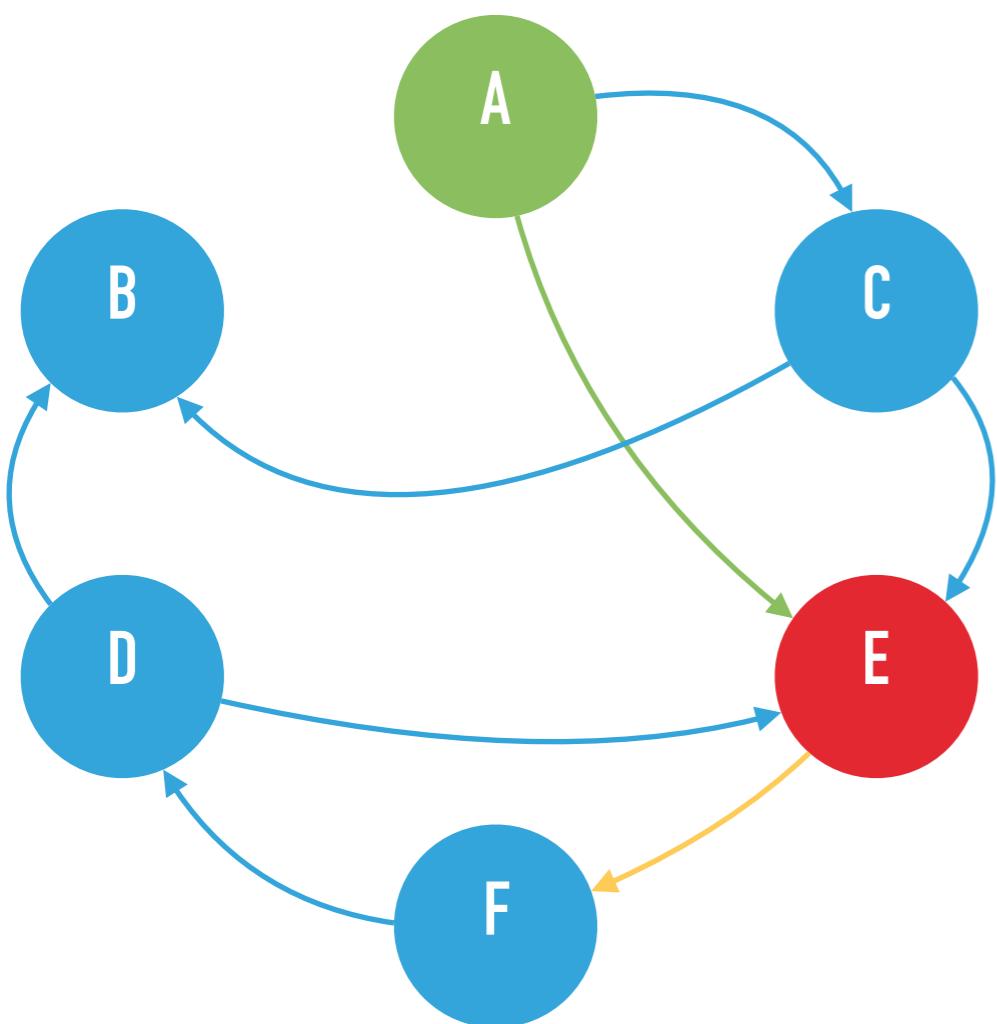
Paths to Consider (Queue)



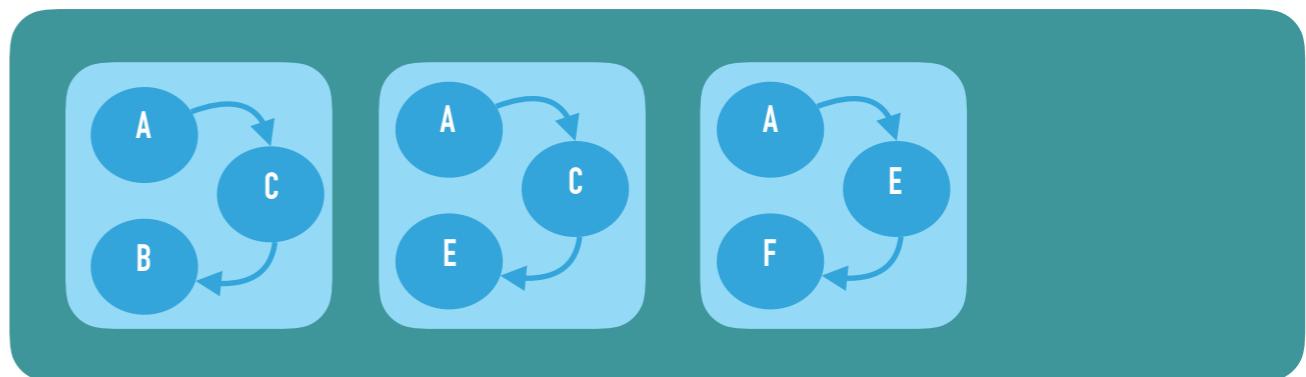
Current Path



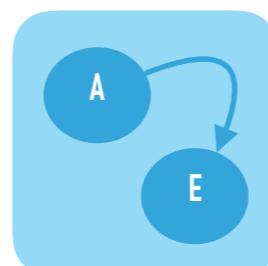
## BREADTH FIRST SEARCH



Paths to Consider (Queue)



Current Path



YOU NEVER CONSIDER A PATH OF  
LENGTH  $K + 1$

UNTIL YOU'VE CONSIDERED ALL PATHS OF  
LENGTH  $K$  OR SHORTER

# COMPARING DFS AND BFS

## COMPARING DFS AND BFS

### DFS

- ▶ create a path with just start node and push onto stack  $s$
- ▶ while  $s$  is not empty:
  - ▶  $p = s.pop()$
  - ▶  $v = \text{last node of } p$
  - ▶ if  $v$  is end node, you're done
  - ▶ mark  $v$  as visited
  - ▶ for each unvisited neighbor:
    - ▶ create new path and append neighbor
    - ▶ push new path onto  $s$

### BFS

- ▶ create a path with just start node and enqueue into queue  $q$
- ▶ while  $q$  is not empty:
  - ▶  $p = q.dequeue()$
  - ▶  $v = \text{last node of } p$
  - ▶ if  $v$  is end node, you're done
  - ▶ mark  $v$  as visited
  - ▶ for each unvisited neighbor:
    - ▶ create new path and append neighbor
    - ▶ enqueue new path into  $q$

## COMPARING DFS AND BFS

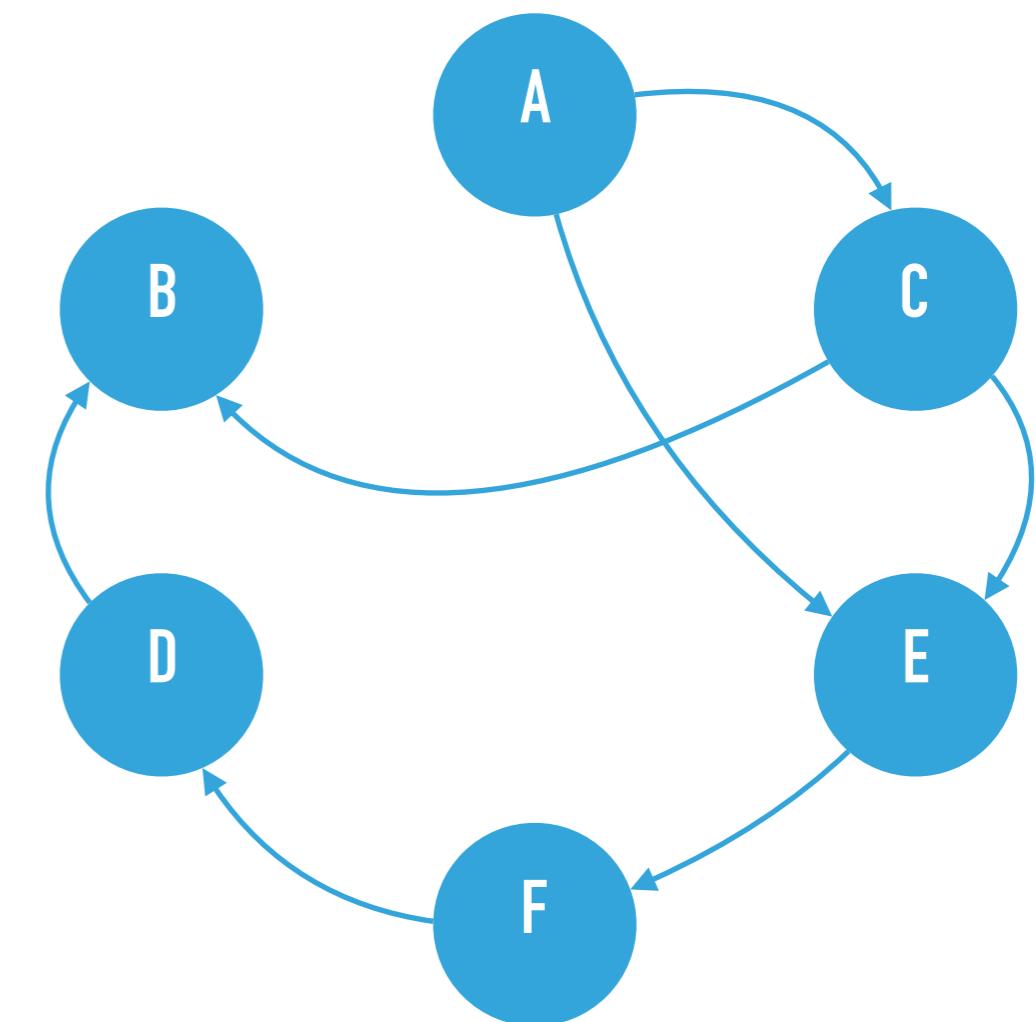
### DFS

- ▶ create a path with just start node and push onto **stack s**
- ▶ while **s** is not empty:
  - ▶ **p = s.pop()**
  - ▶ **v = last node of p**
  - ▶ if **v** is end node, you're done
  - ▶ mark **v** as visited
  - ▶ for each unvisited neighbor:
    - ▶ create new path and append neighbor
    - ▶ **push new path onto s**

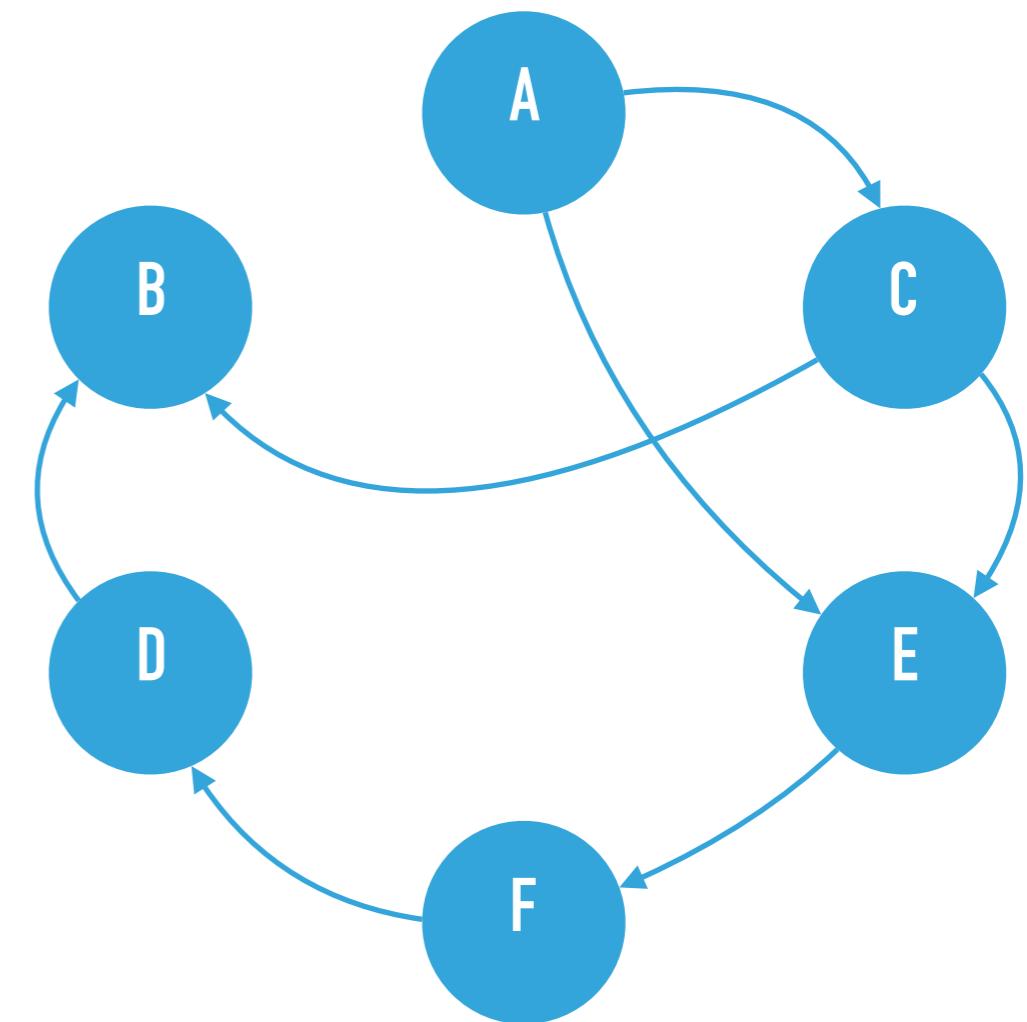
### BFS

- ▶ create a path with just start node and enqueue into **queue q**
- ▶ while **q** is not empty:
  - ▶ **p = q.dequeue()**
  - ▶ **v = last node of p**
  - ▶ if **v** is end node, you're done
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  - ▶ for each unvisited neighbor:
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    - ▶ **enqueue new path into q**

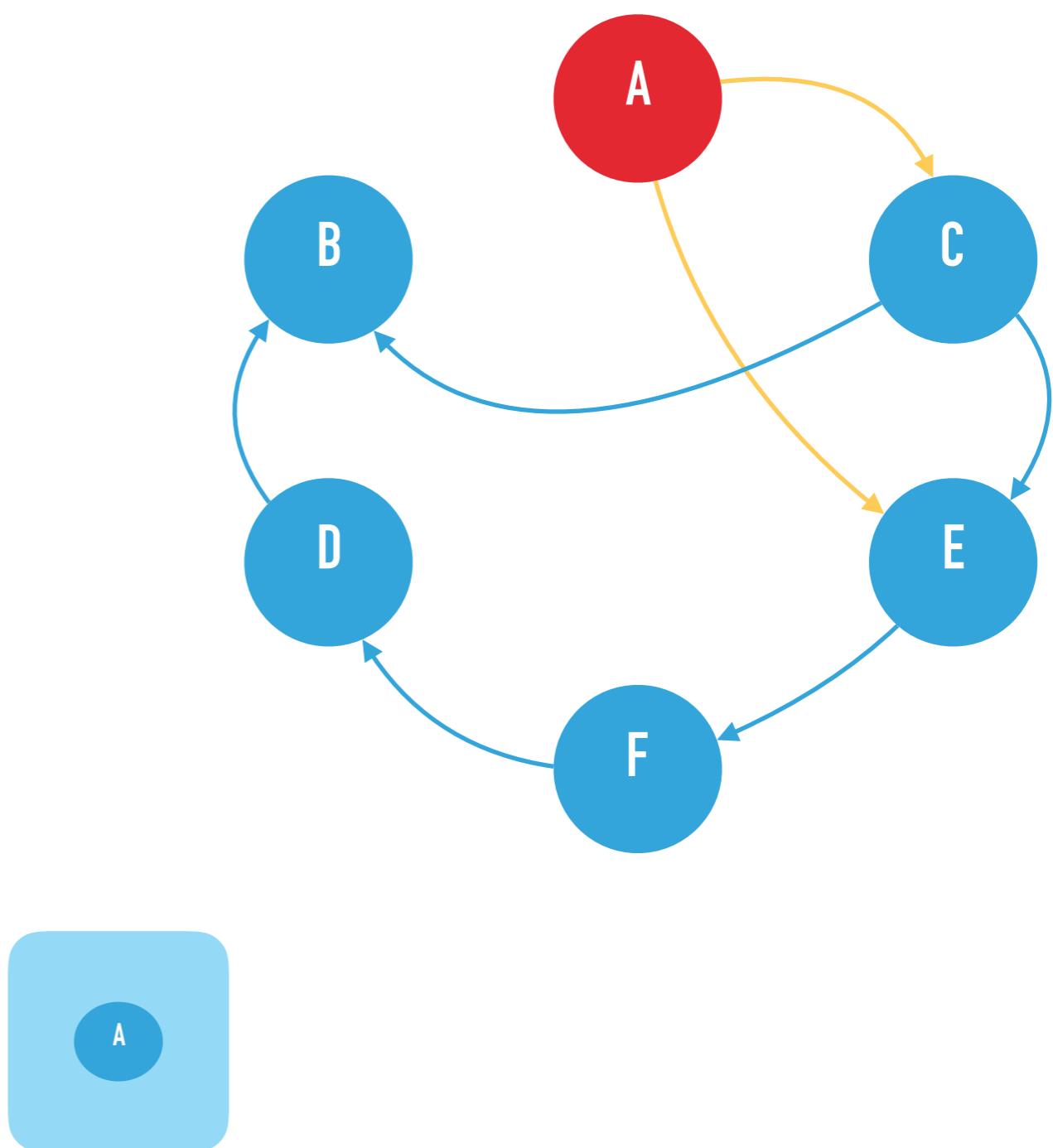
## THE GRAPH SEARCH TO-DO LIST



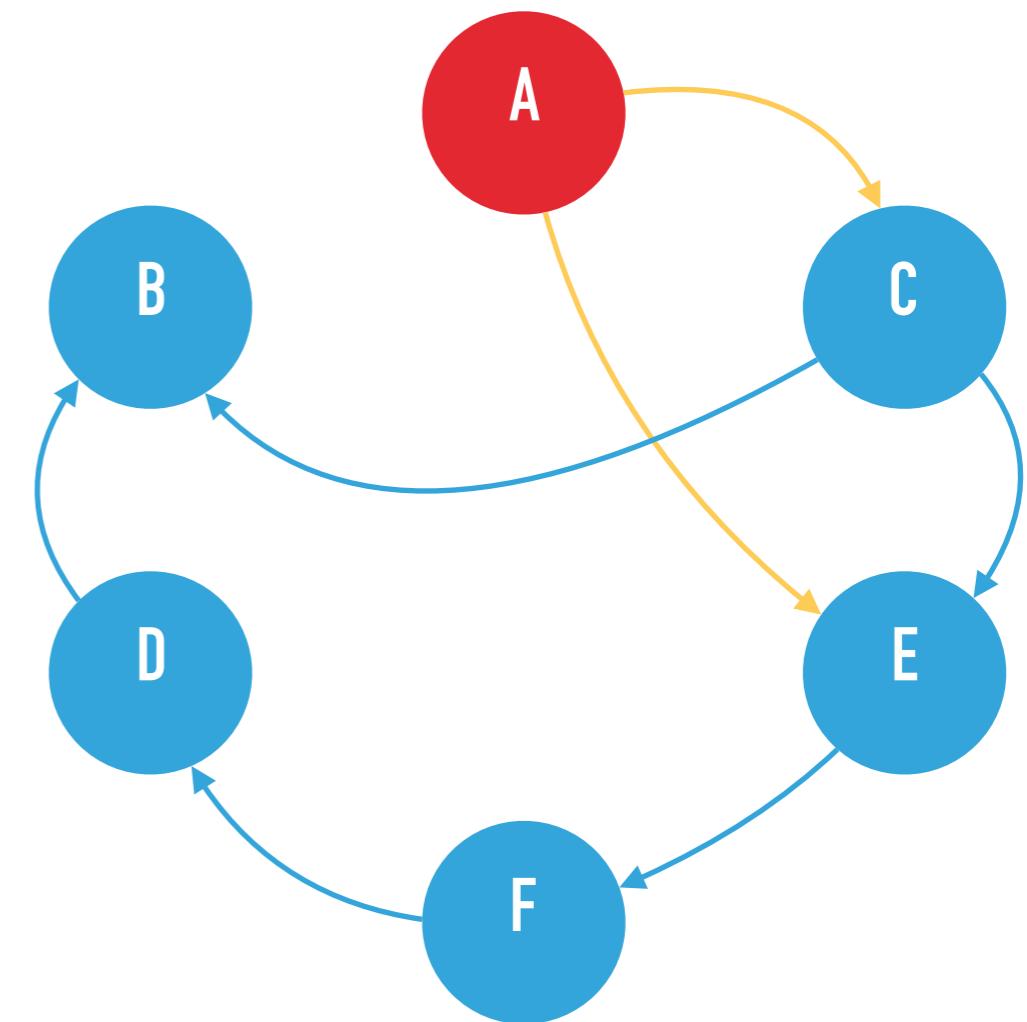
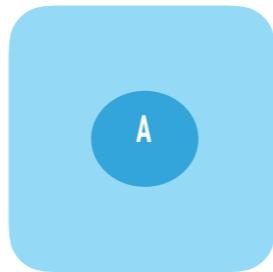
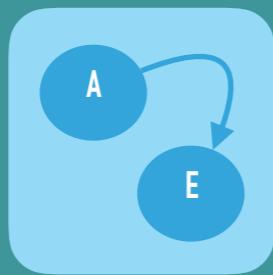
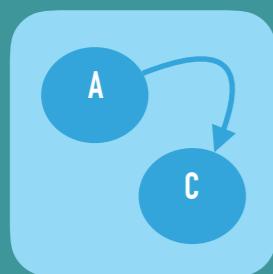
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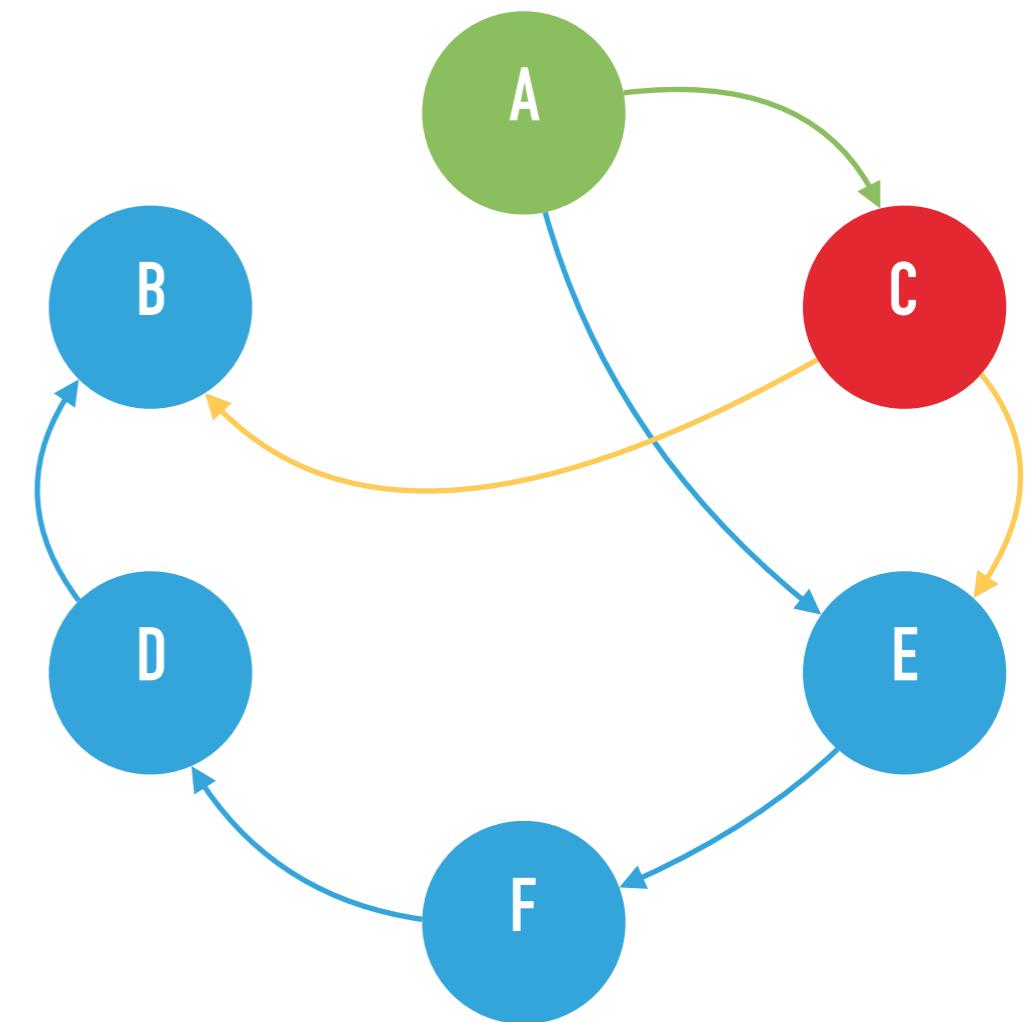
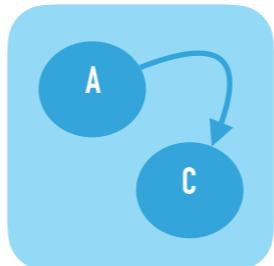
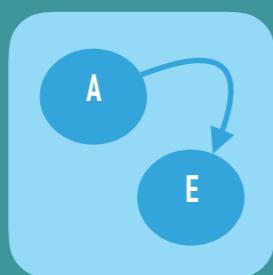
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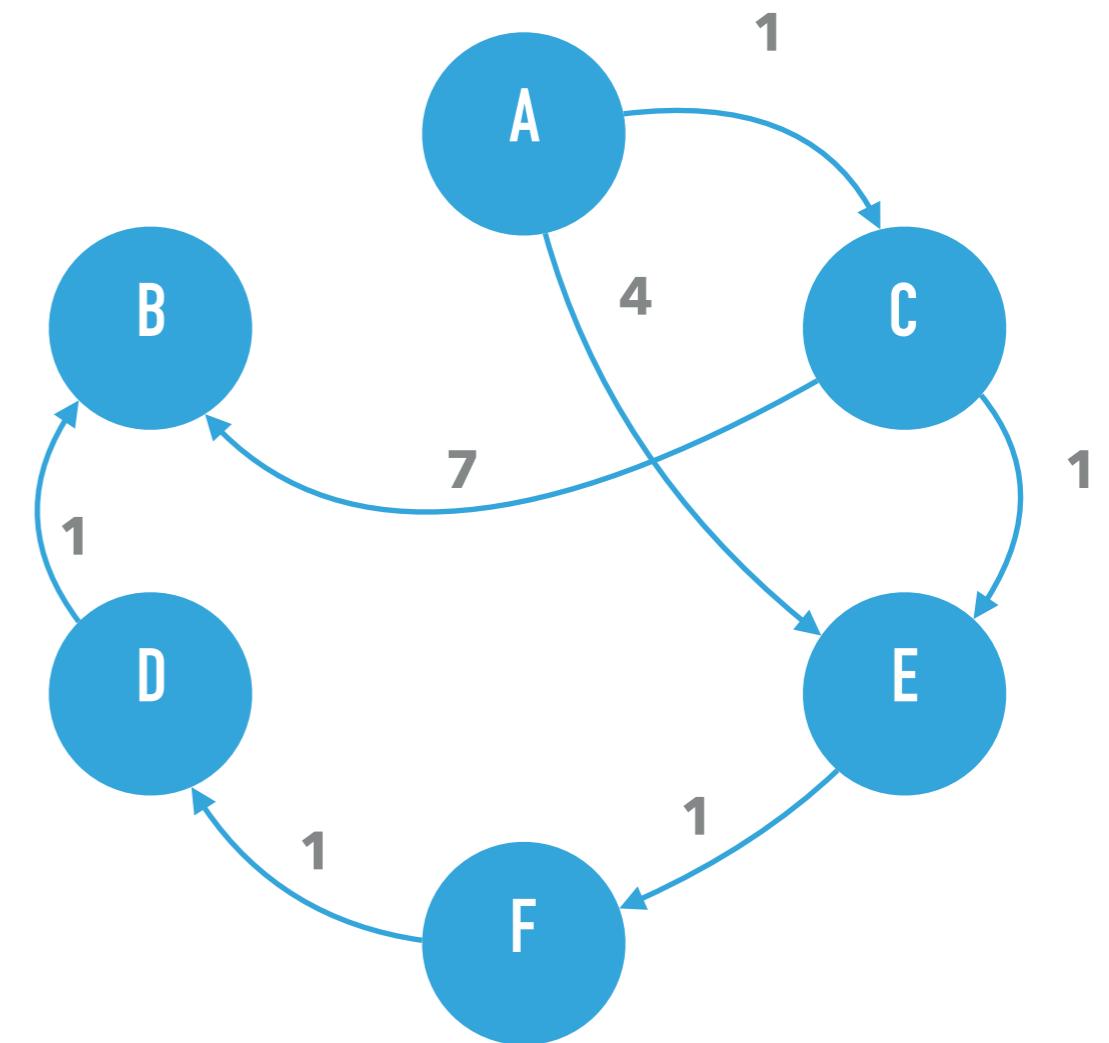


## THE GRAPH SEARCH TO-DO LIST

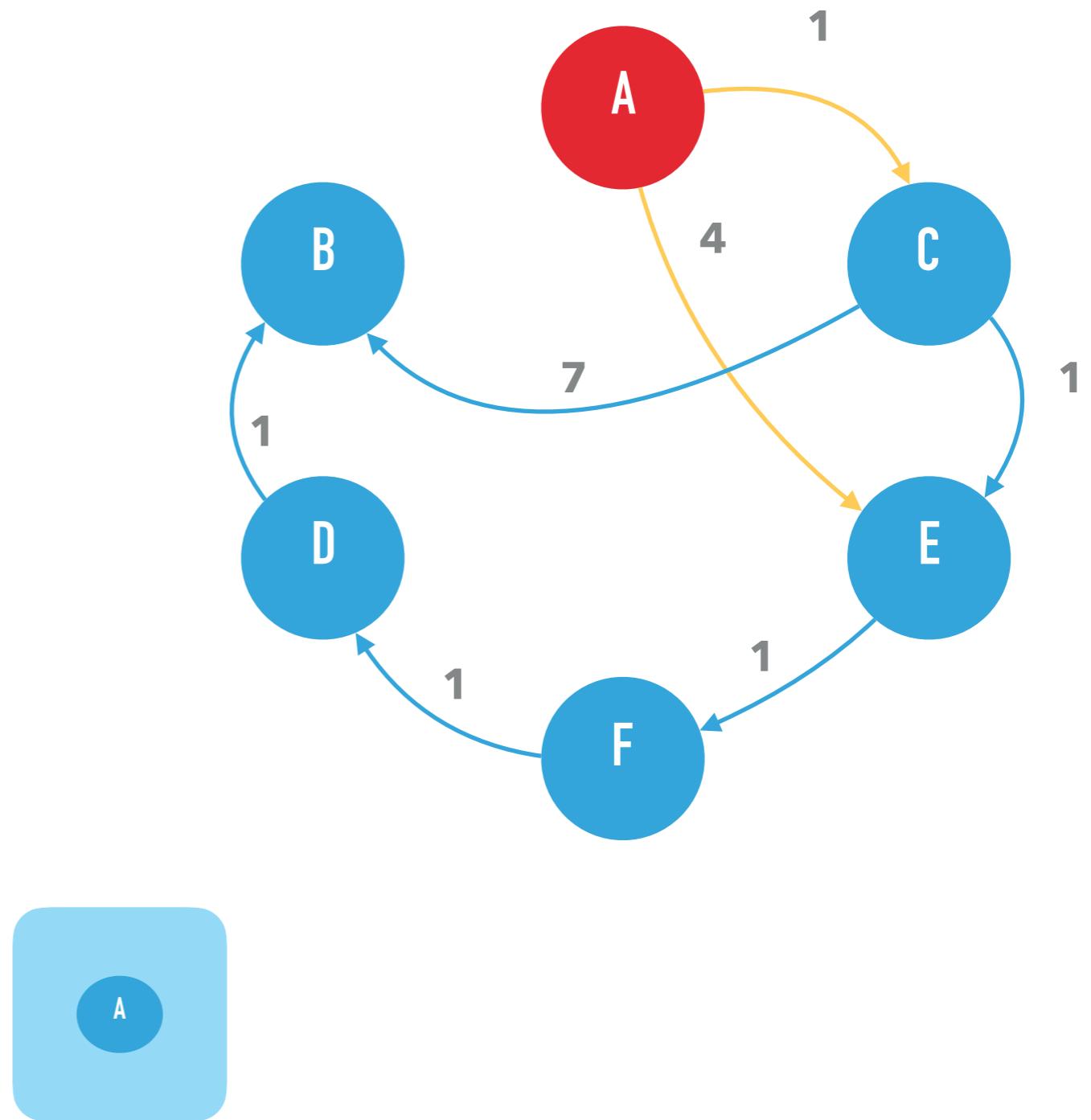


**WEIGHTY  
DECISIONS**

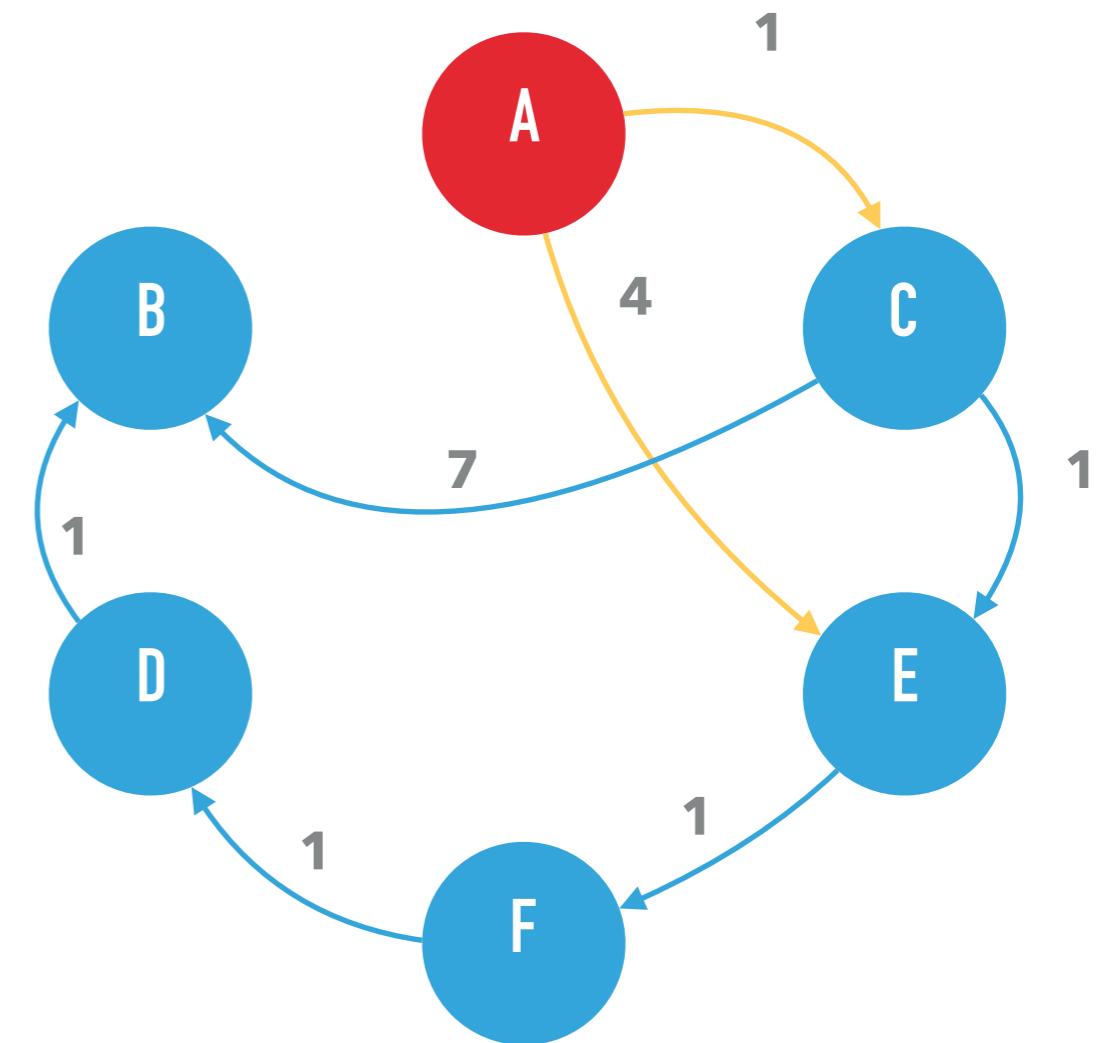
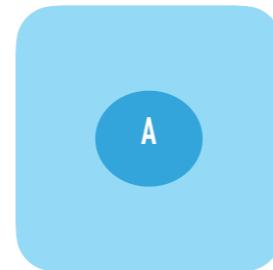
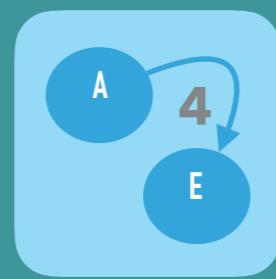
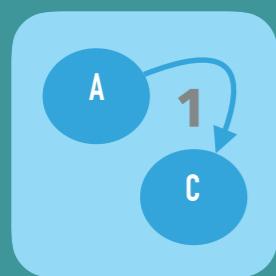
## DEALING WITH WEIGHTY TOPICS



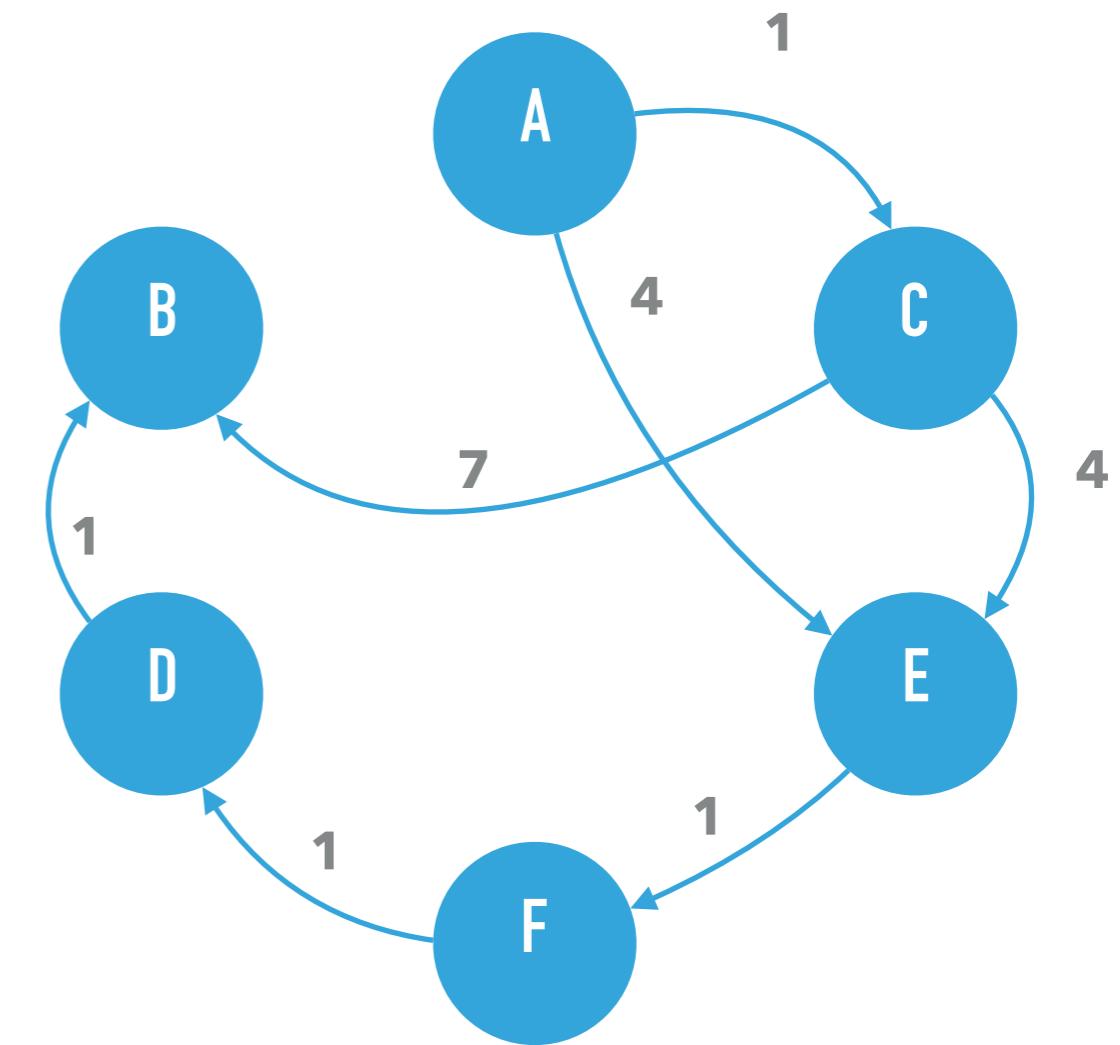
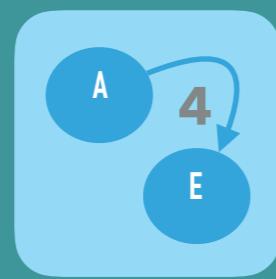
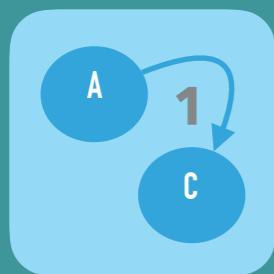
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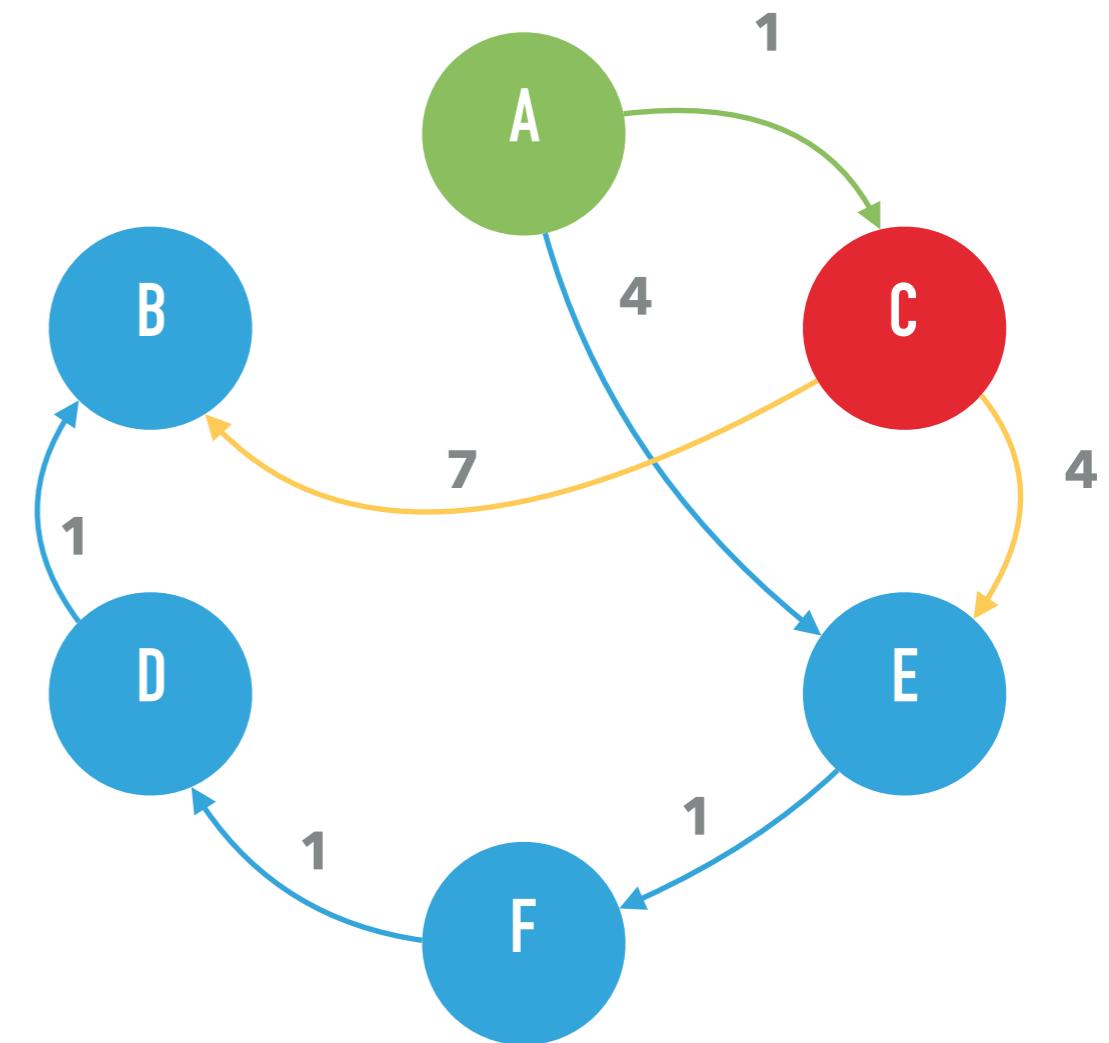
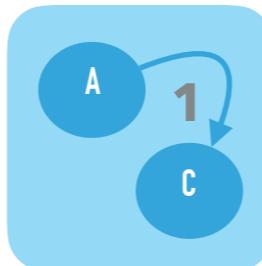
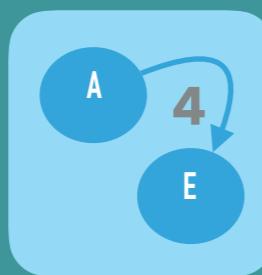
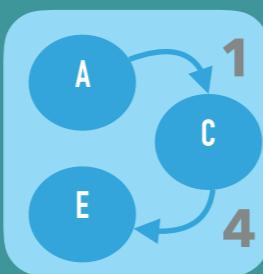
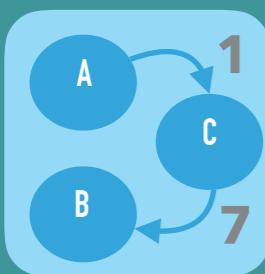
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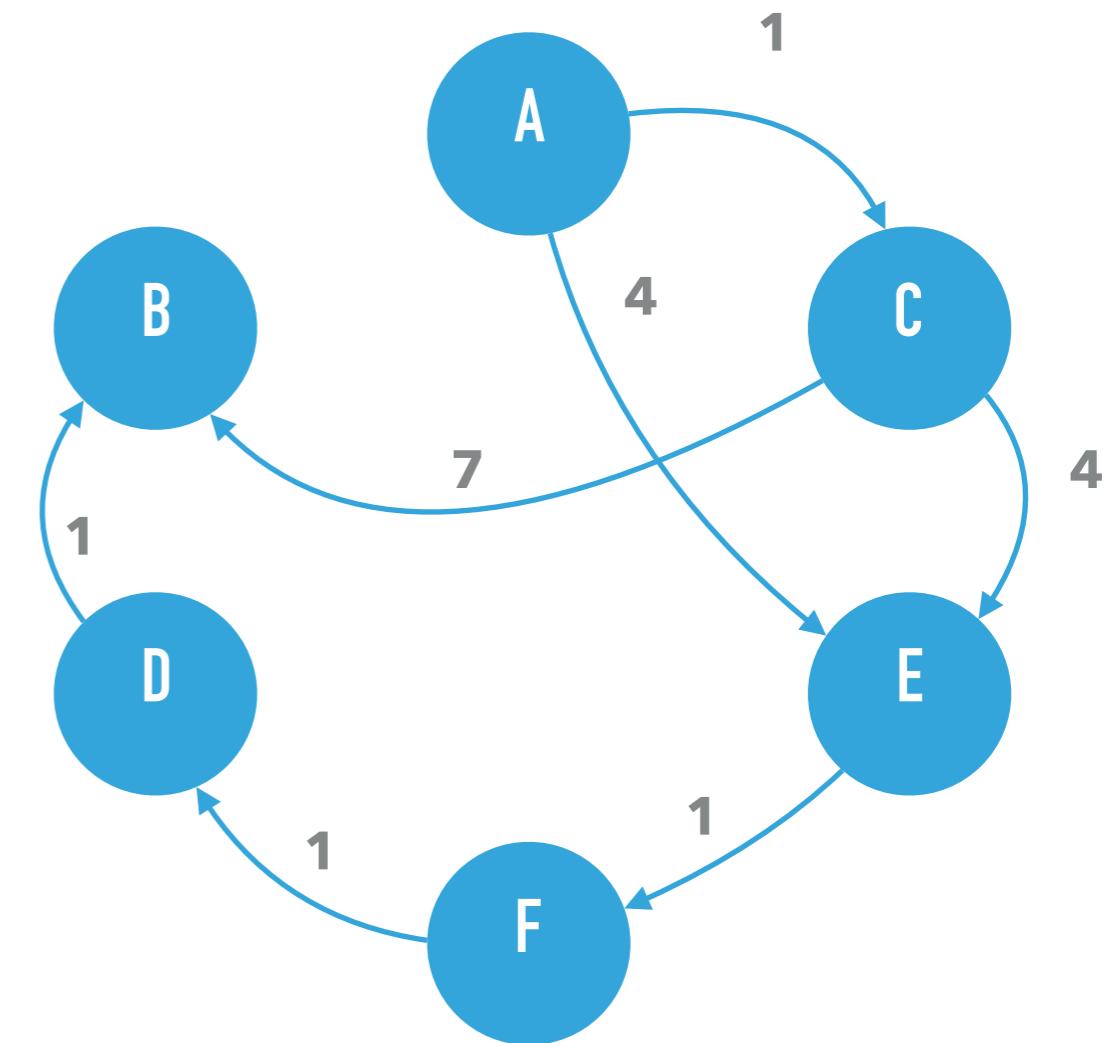
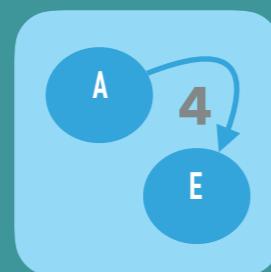
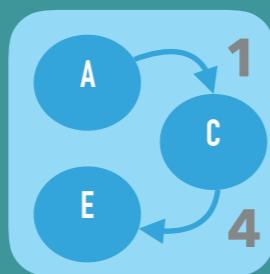
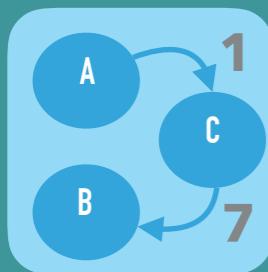
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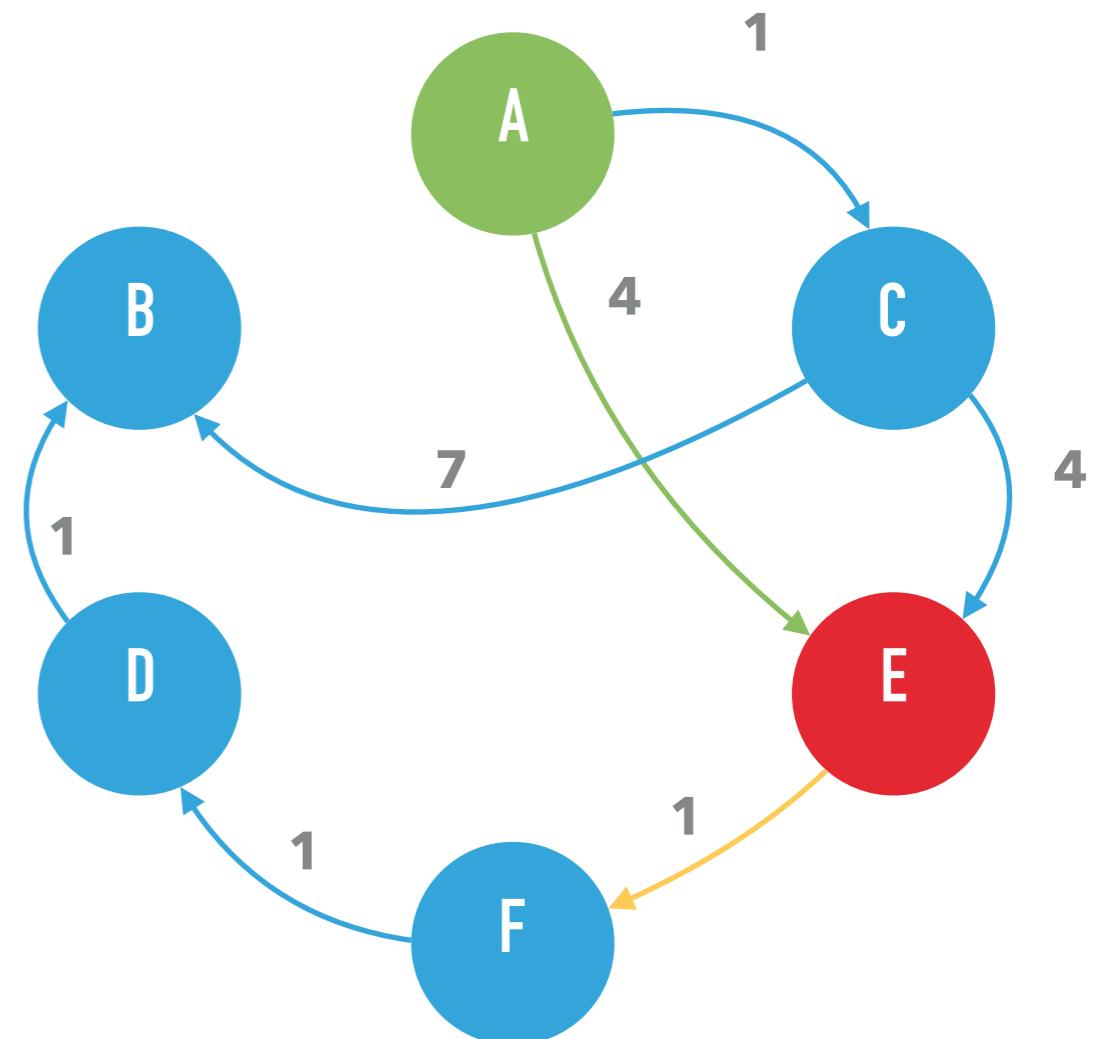
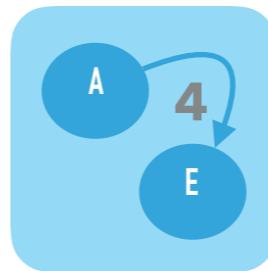
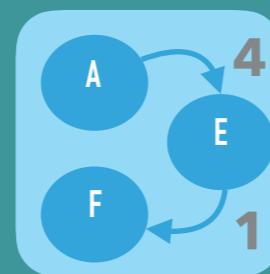
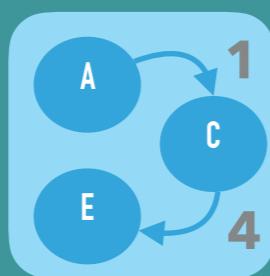
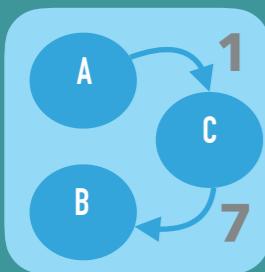
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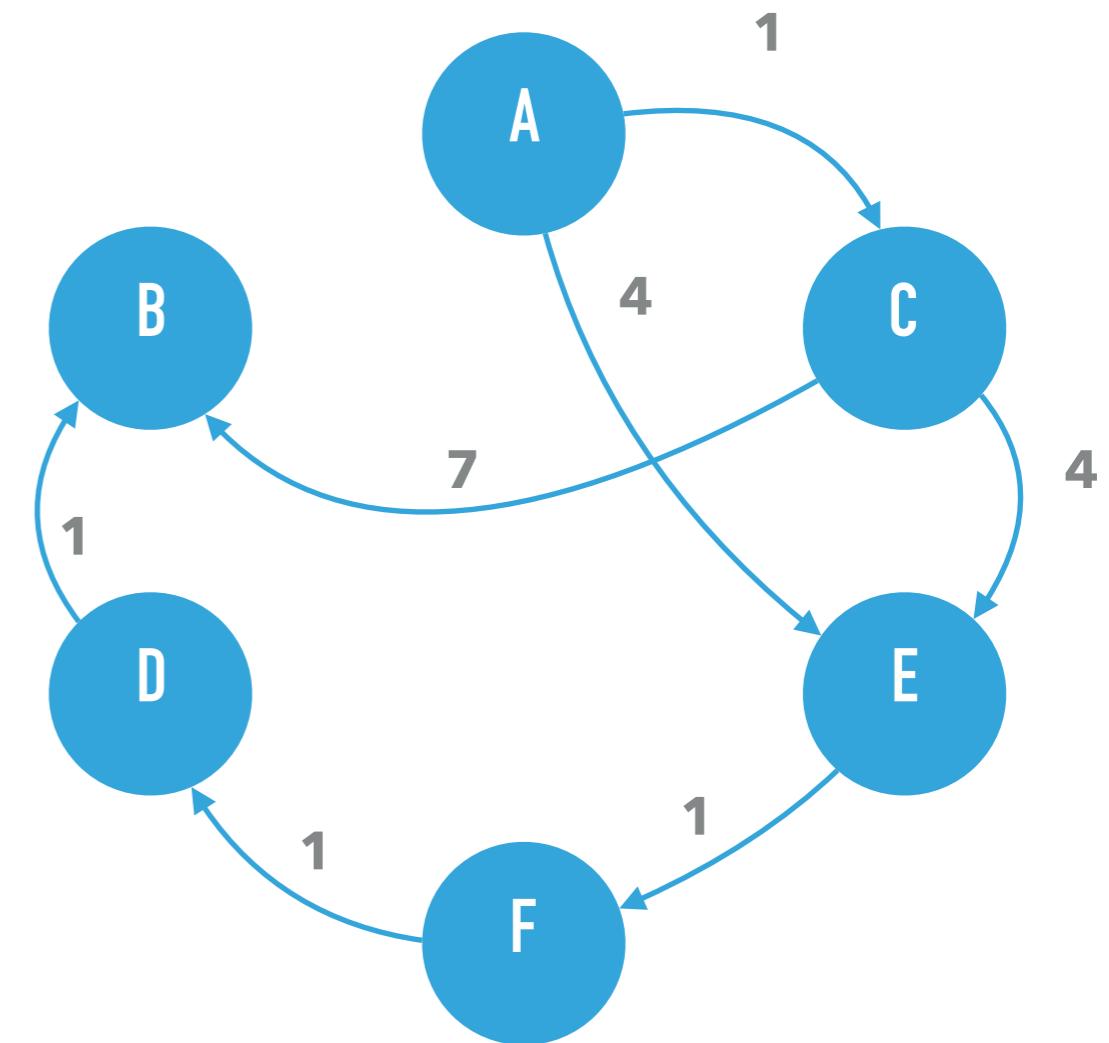
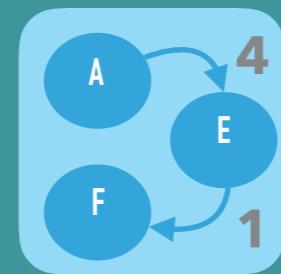
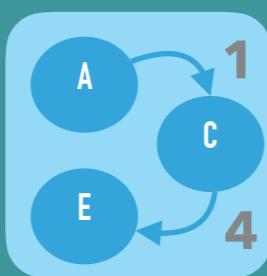
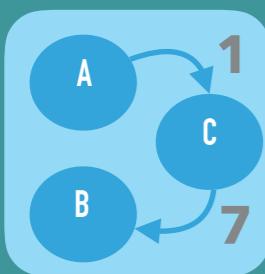
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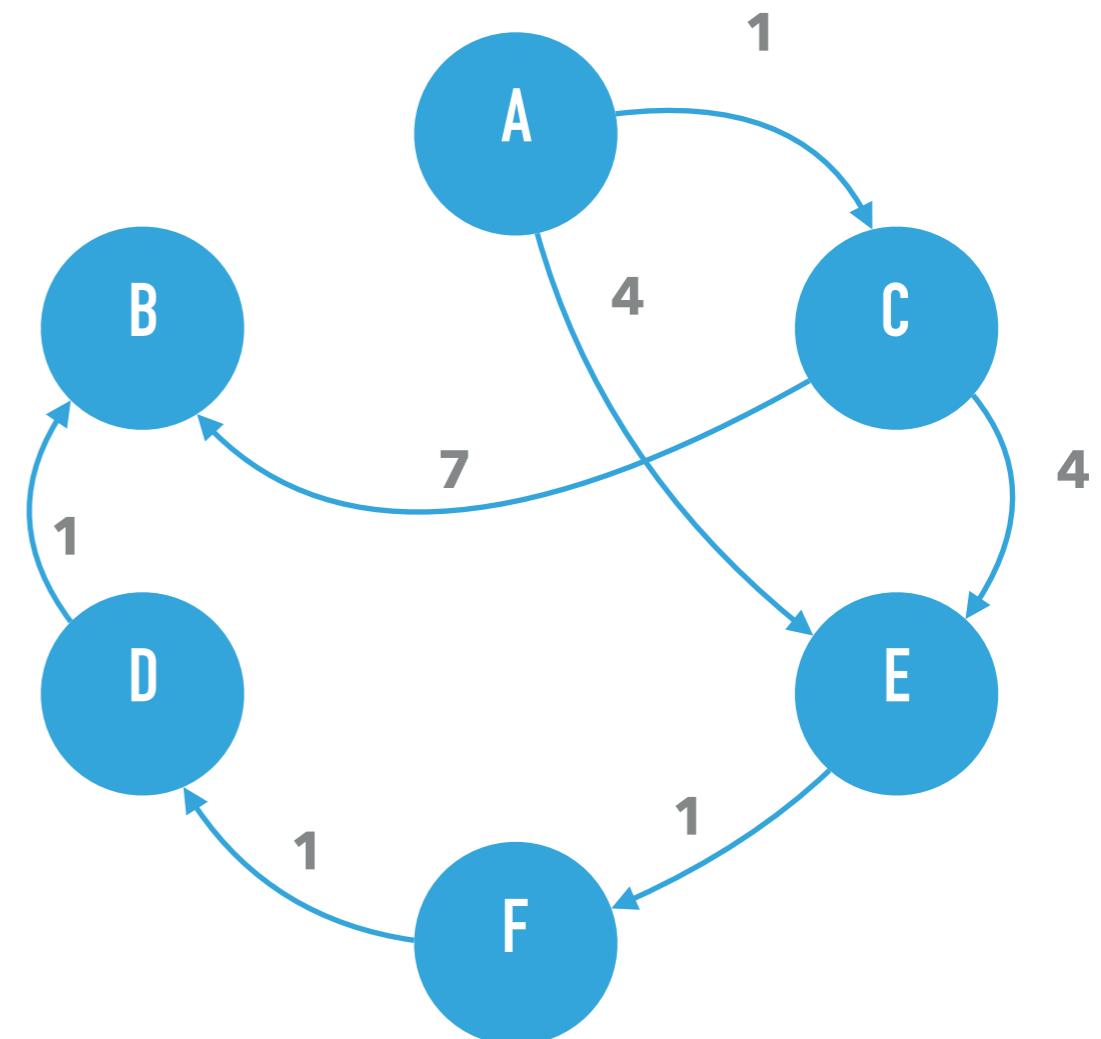
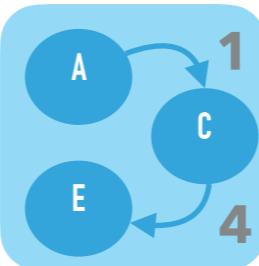
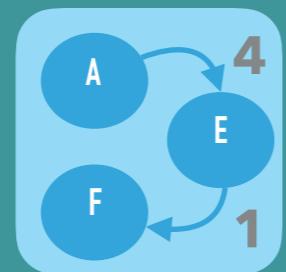
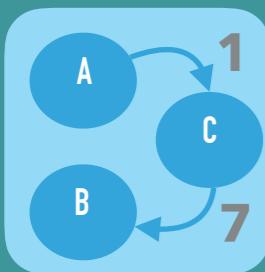
## DEALING WITH WEIGHTY TOPICS



## DEALING WITH WEIGHTY TOPICS



## DEALING WITH WEIGHTY TOPICS



IN DIJKSTRA'S ALGORITHM,

THE TODO LIST IS A PRIORITY  
QUEUE

### DIJKSTRA'S ALGORITHM (PSEUDOCODE)

- ▶ create a path with just start node and enqueue into priority queue q
- ▶ while q is not empty
  - ▶  $p = q.dequeue()$
  - ▶  $v = \text{last node of } p$
  - ▶ if  $v$  is end node, you're done
  - ▶ if you've seen  $v$  before, skip it
  - ▶ mark  $v$  as visited
  - ▶ for each unvisited neighbor:
    - ▶ create new path and append neighbor
    - ▶ enqueue new path into q

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  - ▶ for each unvisited neighbor:
    - ▶ create new path and append neighbor
    - ▶ enqueue new path into q with priority  $\text{pathLength}$

# DIJKSTRA'S ODDS AND ENDS

- ▶ **create a path with just start node and enqueue into priority queue q**
- ▶ while q is not empty
  - ▶  $p = q.dequeue()$
  - ▶  $v = \text{last node of } p$
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- ▶ What do you initialize the weight of the path to?

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  - ▶ mark `v` as visited
  - ▶ for each unvisited neighbor:
    - ▶ create new path and append neighbor
    - ▶ enqueue new path into q with priority `pathLength`
- ▶ What do you initialize the weight of the path to?
  - ▶ Zero should be fine

### DIJKSTRA'S ODDS AND ENDS

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    - ▶ create new path and append neighbor
    - ▶ enqueue new path into q with priority  $\text{pathLength}$
- ▶ Can't I just return the path as soon as I find the end node? Why wait until I dequeue?
  - ▶ This is one of the most common mistakes people make with Dijkstra's!
  - ▶ It's possible a path with a lower priority gets enqueued in the meantime.

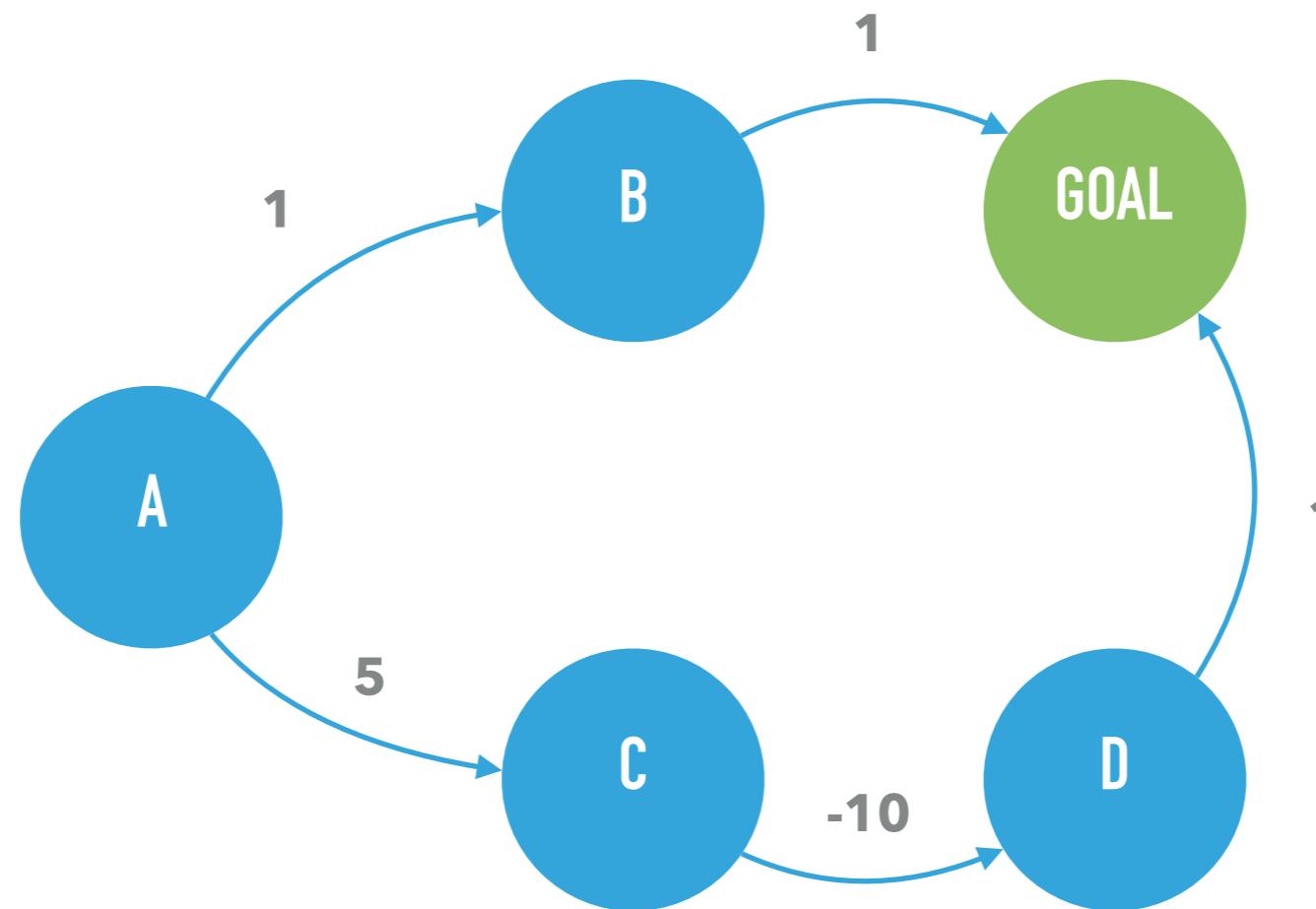
## DIJKSTRA'S ODDS AND ENDS

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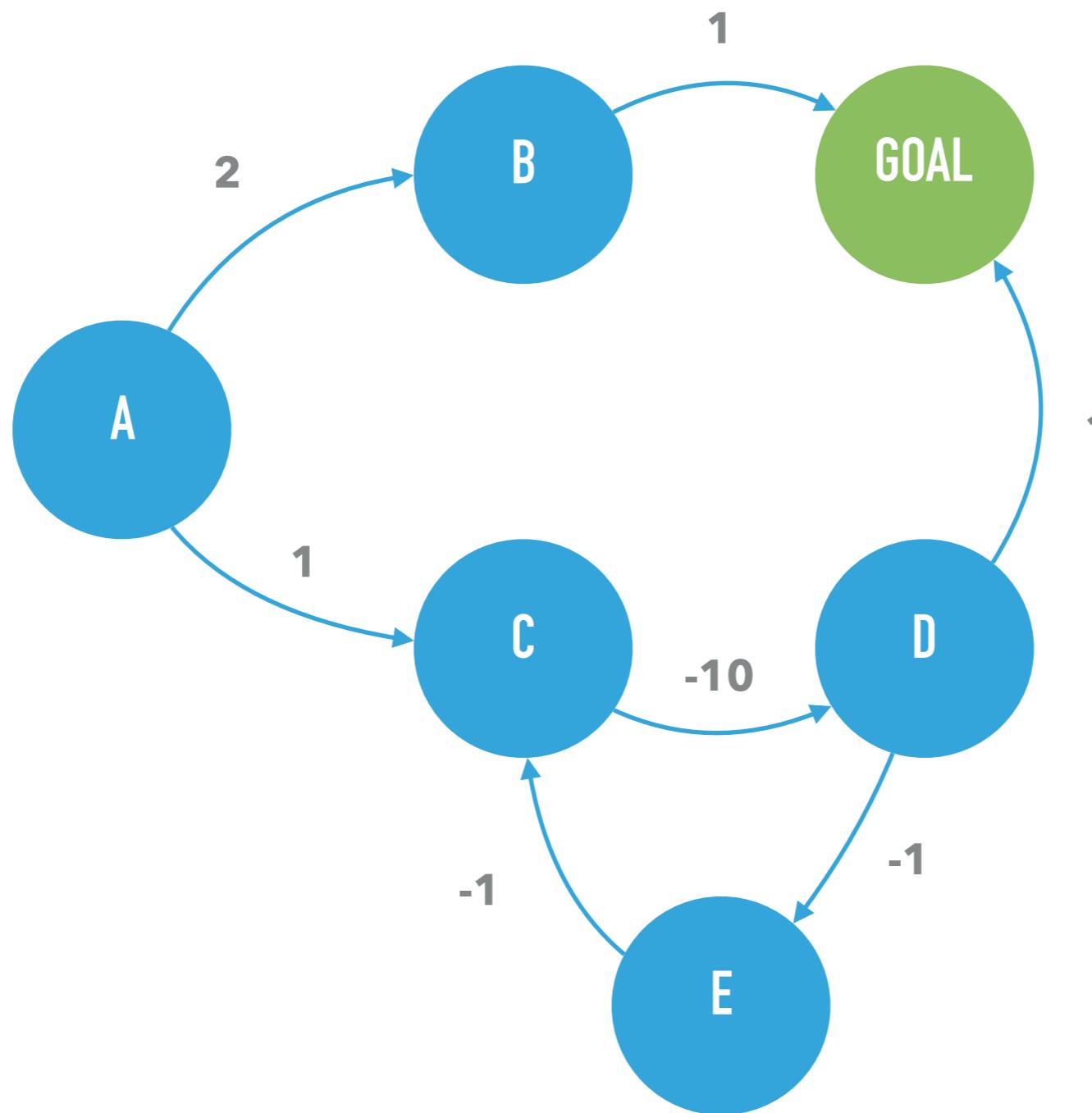
## DIJKSTRA'S ODDS AND ENDS

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  - ▶ for each unvisited neighbor:
    - ▶ create new path and append neighbor
    - ▶ enqueue new path into q with priority `pathLength`
- ▶ Why would you skip the node just because you've seen it before?
  - ▶ If you've seen the node before, that means you've already found a shorter path to it.
  - ▶ Any path that follows from this one already has a shorter equivalent
  - ▶ **The first path you find to `v` will be the shortest path to `v`**

## NEGATIVE EDGES



## NEGATIVE CYCLES









1

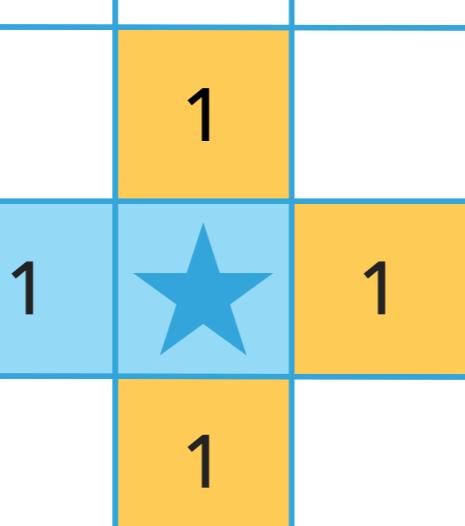
1



1

1













			2	
		2	1	2
2	1	★	1	
	2	1	2	
		2		

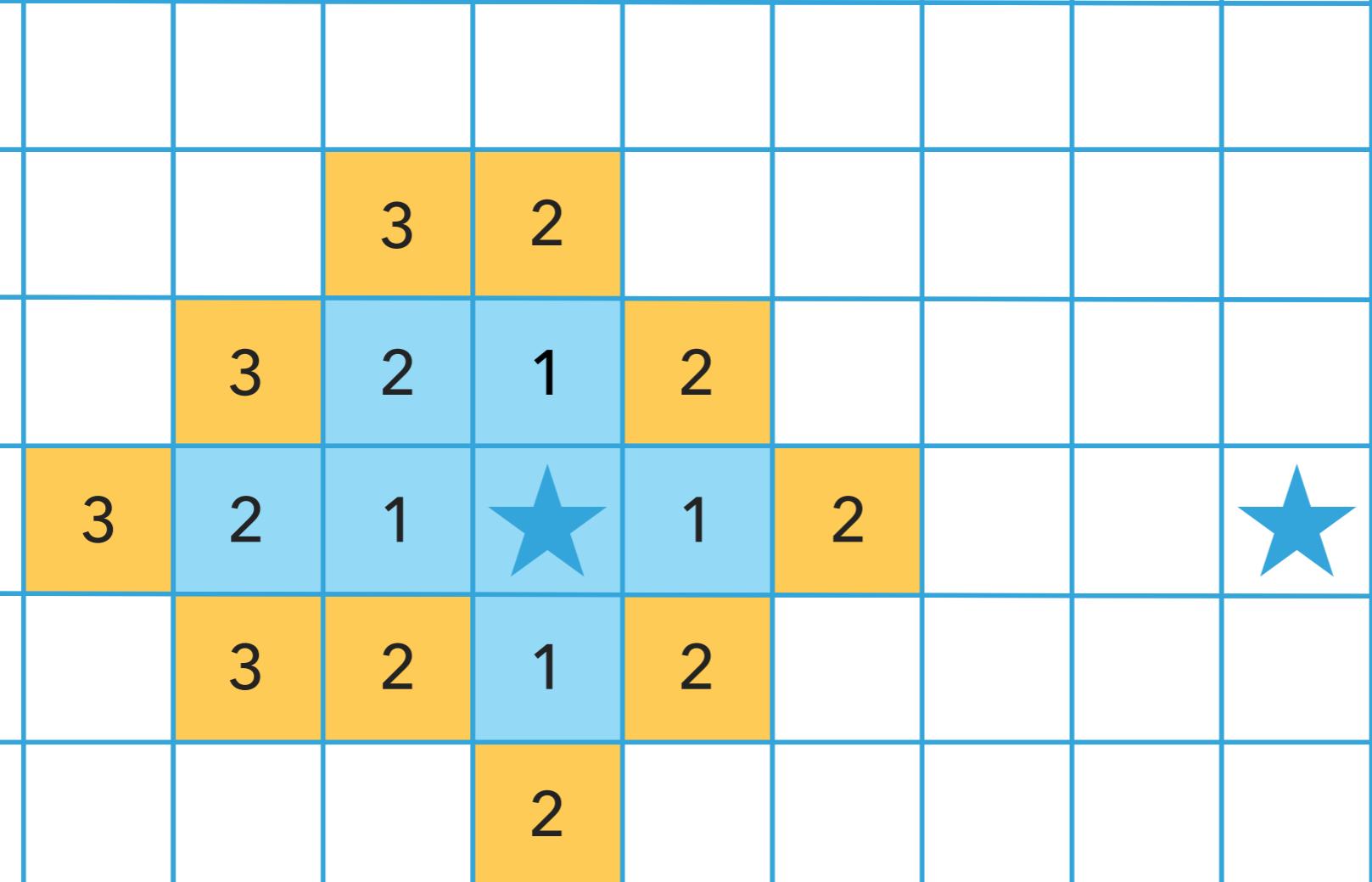
			2	
		2	1	2
2	1	★	1	★
	2	1	2	
		2		

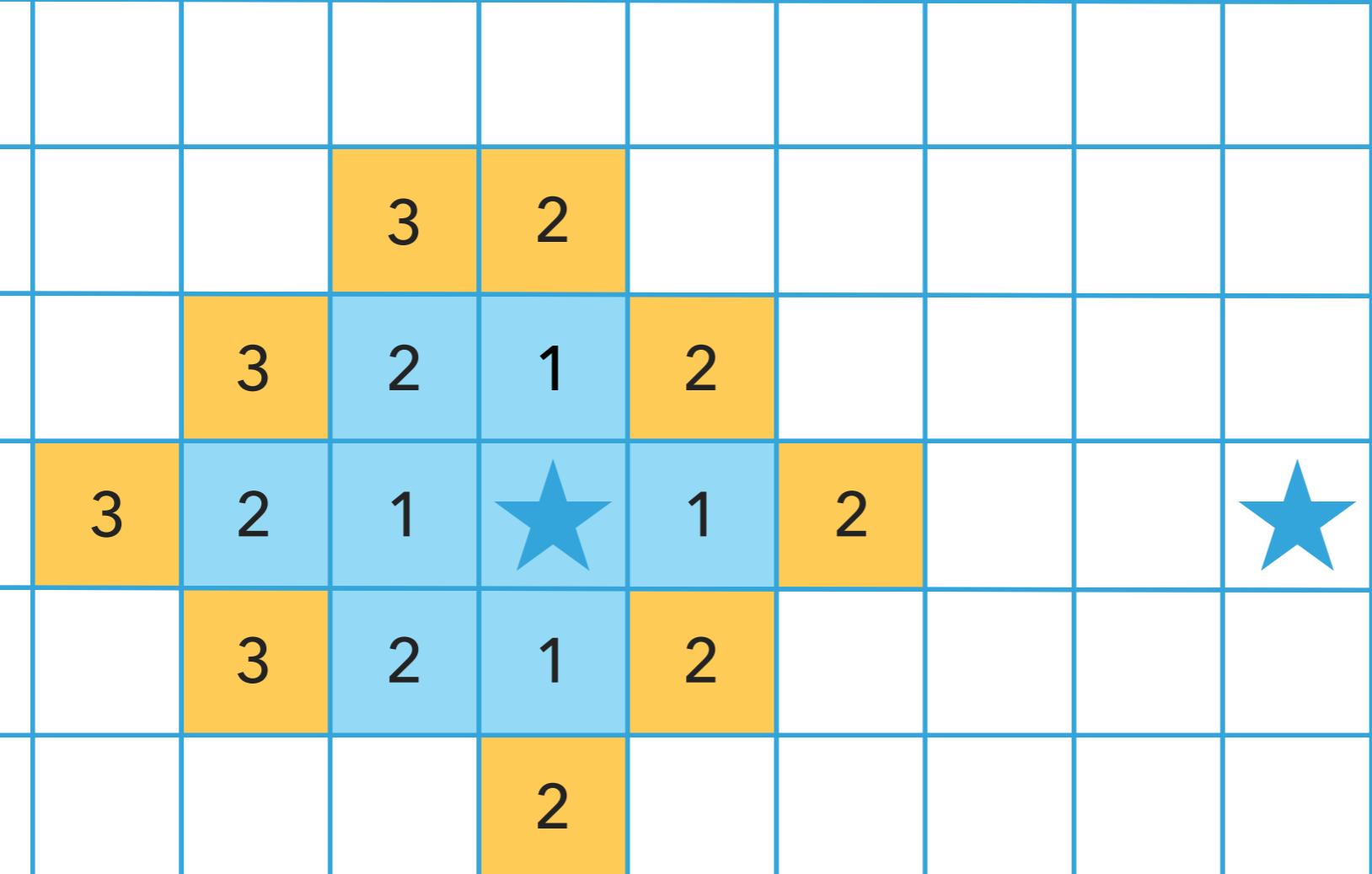
			2					
		2	1	2				
2	1	★	1	2			★	
	2	1	2					
		2						

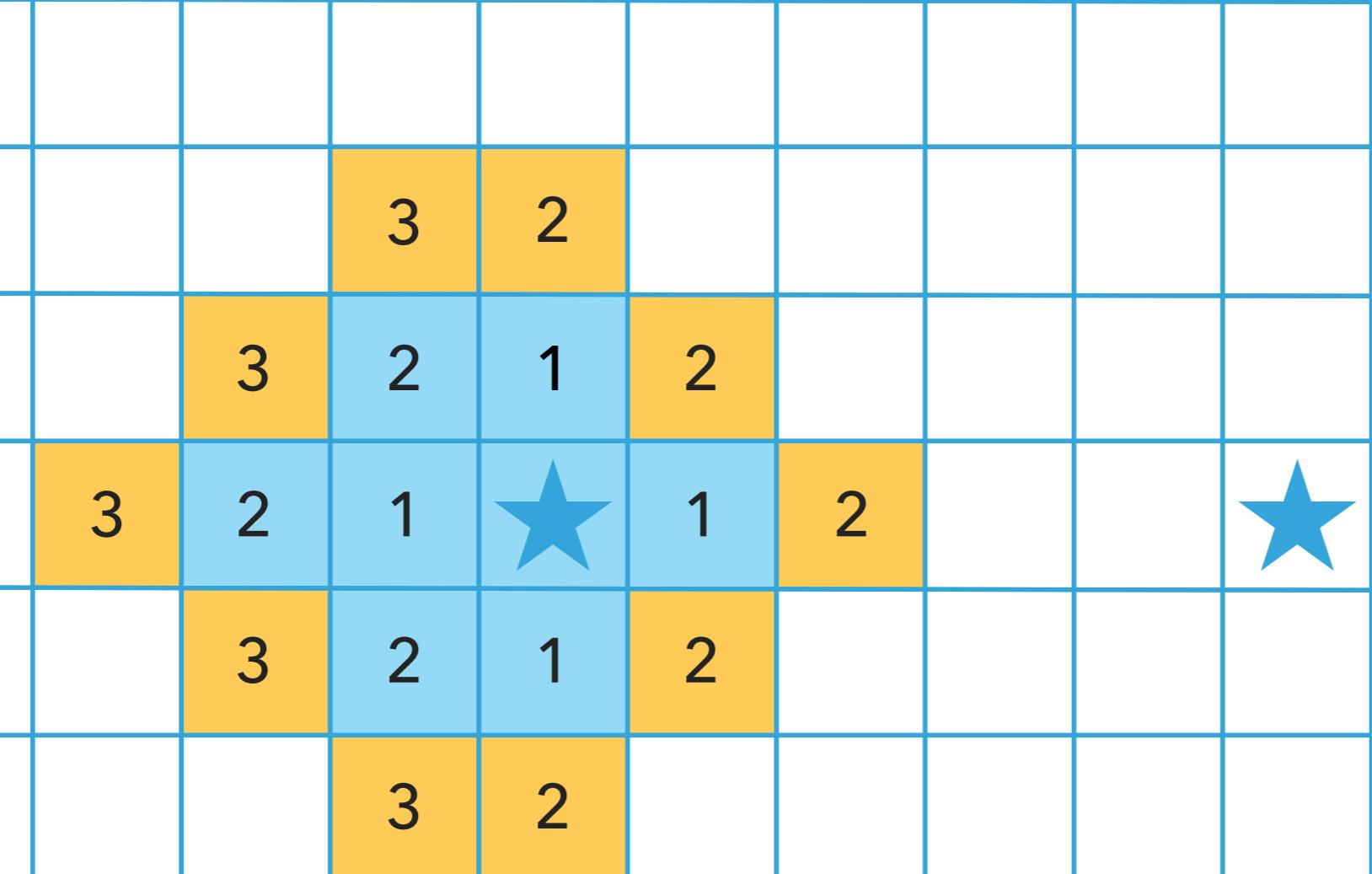
			2						
		2	1	2					
2	1	★	1	2			★		
	2	1	2						
		2							

				2			
		3	2	1	2		
3	2	1	★	1	2		★
	3	2	1	2			
			2				

				2			
		3	2	1	2		
	3	2	1	★	1	2	★
	3	2	1	2			
			2				









				3			
				3	2	3	
				3	2	1	2
	3	2	1	★	1	2	
	3	2	1	2			★
		3	2				

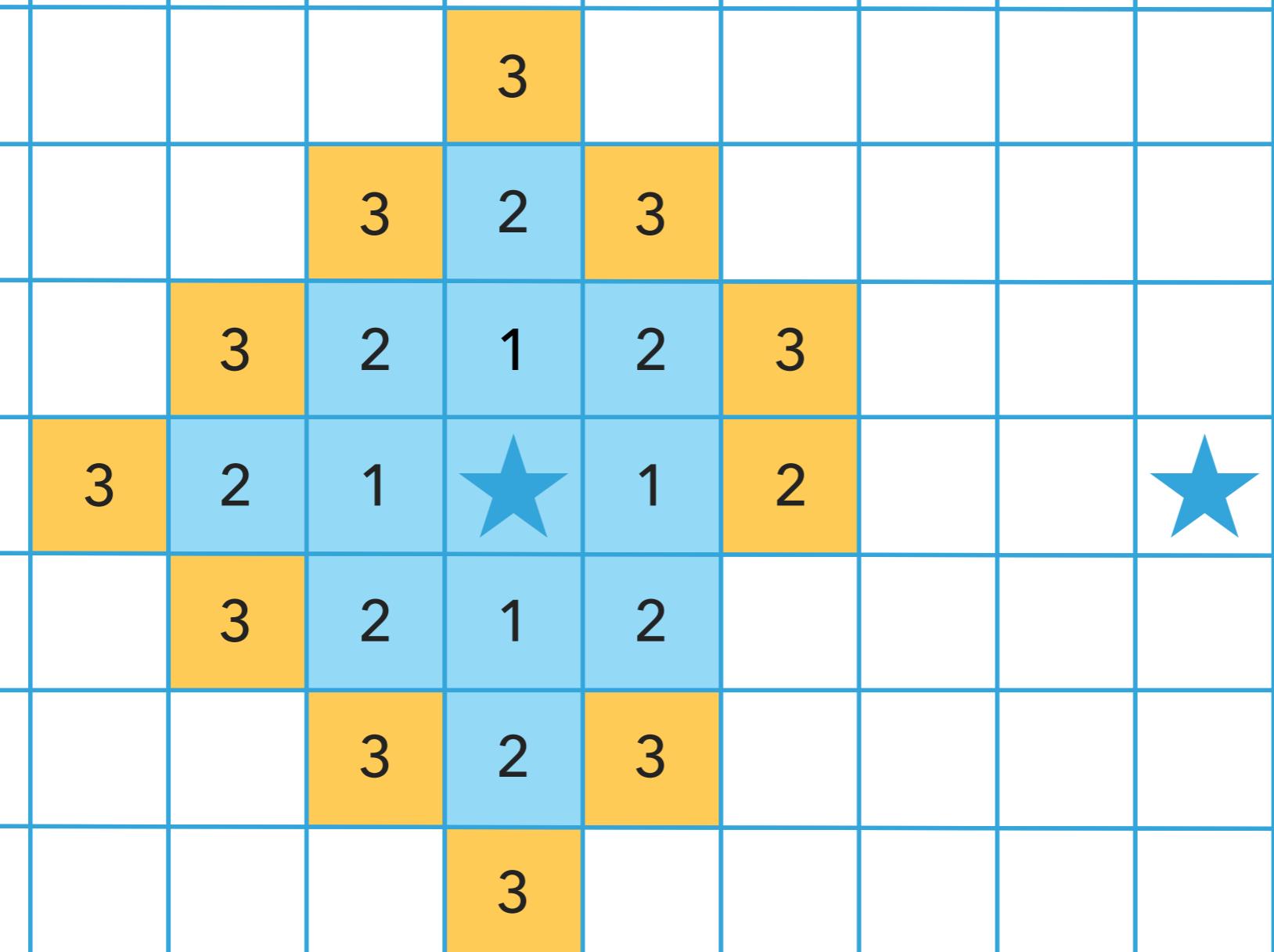
				3			
				3	2	3	
				3	2	1	2
	3	2	1	★	1	2	
	3	2	1	2			★
		3	2				

				3			
			3	2	3		
		3	2	1	2		
	3	2	1	★	1	2	
	3	2	1	2			★
		3	2	3			
			3				

				3			
			3	2	3		
		3	2	1	2		
	3	2	1	★	1	2	
	3	2	1	2			★
		3	2	3			
			3				

				3			
				3	2	3	
				3	2	1	2
				3	2	1	2
				3	2	1	2
				3	2	3	
				3			





			3			
		3	2	3		
	3	2	1	2	3	
3	2	1	★	1	2	
	3	2	1	2	3	
	3	2	3			
		3				



			3				
		3	2	3			
	3	2	1	2	3		
3	2	1	★	1	2		★
	3	2	1	2	3		
	3	2	3				
		3					

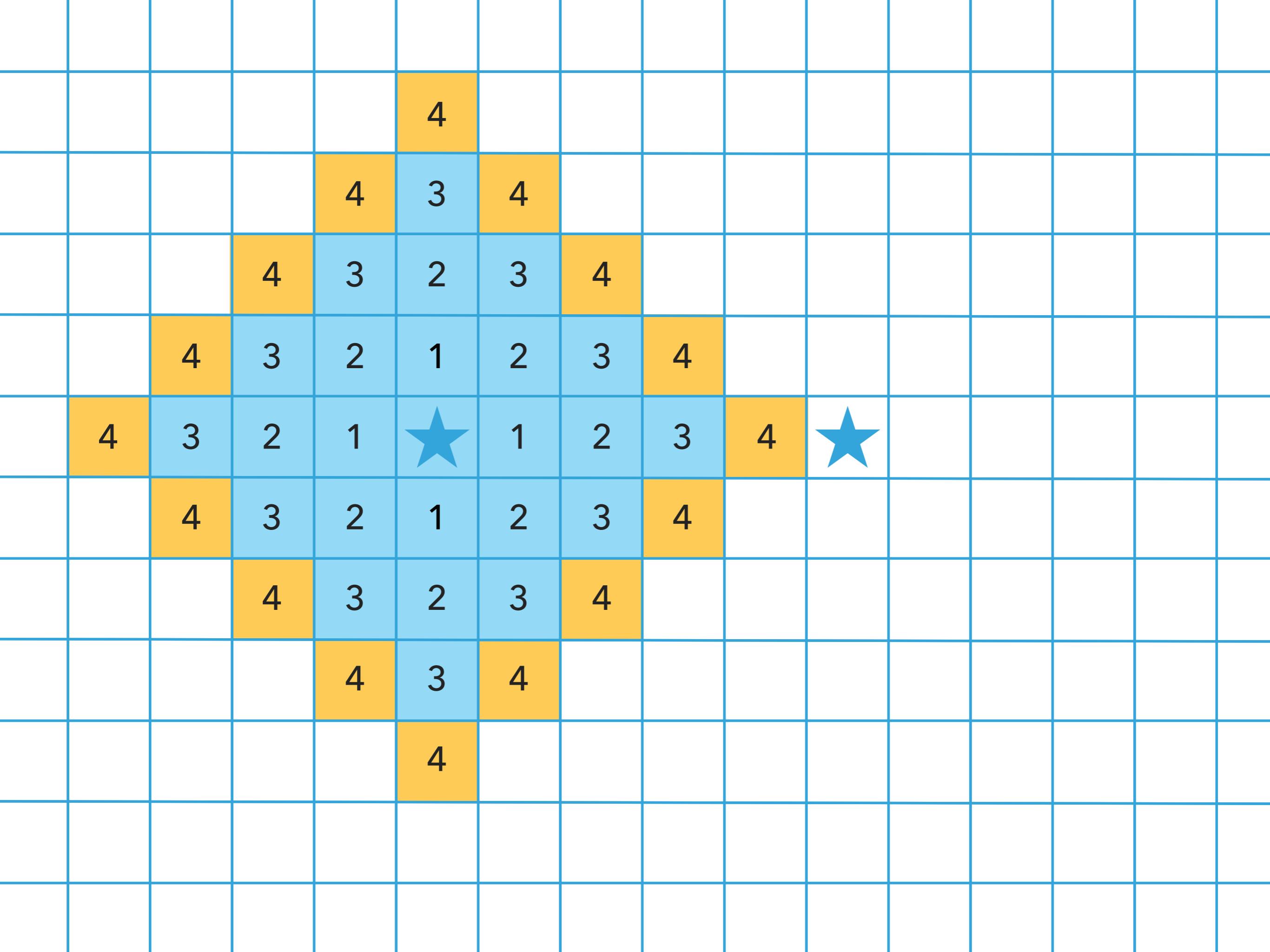

A 9x9 grid puzzle with numbered cells and two star markers. The grid contains the following values:



The values in the grid are as follows:

- Row 1: 3, 3, 2, 3, 3, 3, 3, 3, 3
- Row 2: 3, 3, 2, 1, 2, 3, 3, 3, 3
- Row 3: 3, 2, 1, 3, 1, 2, 3, 3, 3
- Row 4: 3, 2, 1, 2, 3, 3, 3, 3, 3
- Row 5: 3, 2, 3, 3, 3, 3, 3, 3, 3
- Row 6: 3, 3, 2, 3, 3, 3, 3, 3, 3
- Row 7: 3, 3, 3, 3, 3, 3, 3, 3, 3
- Row 8: 3, 3, 3, 3, 3, 3, 3, 3, 3
- Row 9: 3, 3, 3, 3, 3, 3, 3, 3, 3

Two blue star markers are present in the grid, located at (row 4, column 4) and (row 5, column 8).



				5						
			5	4	5					
			5	4	3	4	5			
			5	4	3	2	3	4	5	
	5	4	3	2	1	2	3	4	5	
5	4	3	2	1	★	1	2	3	4	★
5	4	3	2	1	2	3	4	5		
	5	4	3	2	3	4	5			
		5	4	3	4	5				
			5	4	5					
				5						

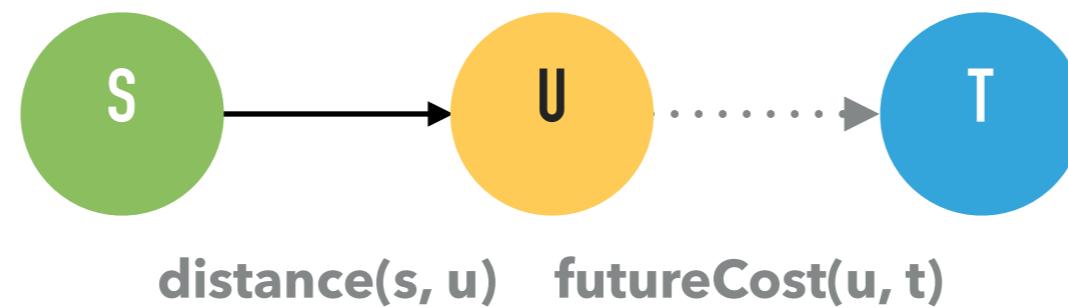
				5						
			5	4	5					
			5	4	3	4	5			
			5	4	3	2	3	4	5	
	5	4	3	2	1	2	3	4	5	
5	4	3	2	1	★	1	2	3	4	★
5	4	3	2	1	2	3	4	5	6	
	5	4	3	2	3	4	5	6		
	5	4	3	4	5	6				
		5	4	5	6					
			5	4	5	6				

DIJKSTRA'S MEASURES THE DISTANCE FROM  
THE START NODE TO THE CURRENT NODE.

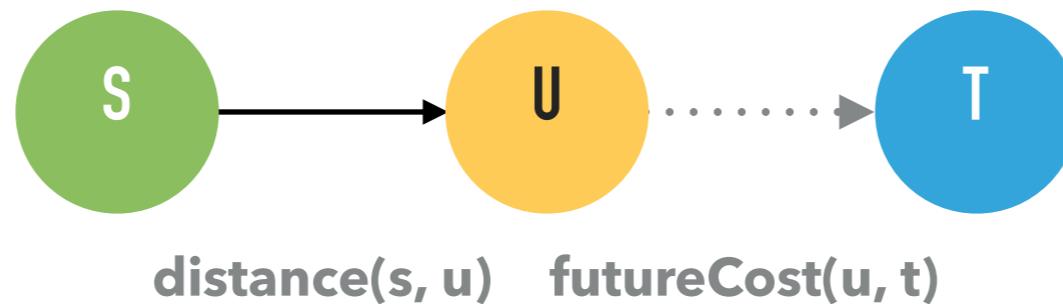
WE WANT THE DISTANCE FROM THE CURRENT  
NODE TO THE DESTINATION.

**SEEING THE  
FUTURE**

## FORMAL DEFINITIONS



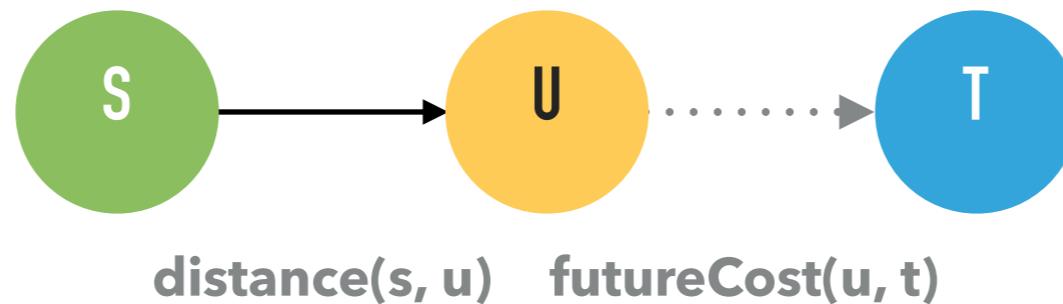
## FORMAL DEFINITIONS



## DIJKSTRA'S

$$priority(u) = distance(s, u)$$

## FORMAL DEFINITIONS

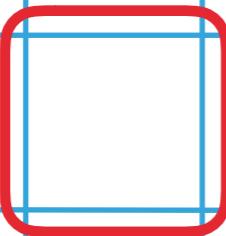


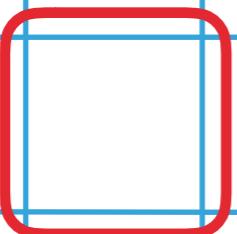
### DIJKSTRA'S

$priority(u) = distance(s, u)$

### IDEAL

$priority(u) = distance(s, u) + futureCost(u, t)$





columns apart

rows apart



```
function futureCost(u, t)  
    return abs(u.row - t.row) + abs(u.col - t.col)
```





1

1



1

1



**1 +  
6**

1



1

1



$1 +$   
**6** $1 +$   
**6****1****1**

$$\begin{array}{ccc} 1 + 6 & & \\ 1 + 6 & \star & 1 \\ 1 + 6 & & \end{array}$$



$1 + 6$

$1 + 6$

$1 + 6$



$1 + 4$



$$\begin{matrix} 1 + \\ 6 \end{matrix} \quad \begin{matrix} 1 + \\ 5 \end{matrix}$$

$$\begin{matrix} 1 + \\ 6 \end{matrix}$$

$$\begin{matrix} 1 + \\ 6 \end{matrix} \quad \begin{matrix} 1 + \\ 5 \end{matrix}$$

$$\begin{matrix} 2 + \\ 3 \end{matrix}$$



1



$$\begin{matrix} 1 + \\ 6 \end{matrix} \quad \begin{matrix} 2 + \\ 5 \end{matrix} \quad \begin{matrix} 3 + \\ 4 \end{matrix}$$

$$\begin{matrix} 1 + \\ 6 \end{matrix}$$

$$\begin{matrix} 1 + \\ 6 \end{matrix} \quad \begin{matrix} 2 + \\ 5 \end{matrix} \quad \begin{matrix} 3 + \\ 4 \end{matrix}$$



1

2

$$\begin{matrix} 3 + \\ 2 \end{matrix}$$



$1 +$ <b>6</b>	$2 +$ <b>5</b>	$3 +$ <b>4</b>	$4 +$ <b>3</b>		
$1 +$ <b>6</b>		1	2	3	$4 +$ <b>1</b>
$1 +$ <b>6</b>	$2 +$ <b>5</b>	$3 +$ <b>4</b>	$4 +$ <b>3</b>		

$$\begin{array}{ccccc} 1+ & 2+ & 3+ & 4+ & 5+ \\ 6 & 5 & 4 & 3 & 2 \end{array}$$

$$\begin{array}{cccccc} 1+ & \star & 1 & 2 & 3 & 4 & 5+ \\ 6 & & & & & & 0 \end{array}$$

$$\begin{array}{ccccc} 1+ & 2+ & 3+ & 4+ & 5+ \\ 6 & 5 & 4 & 3 & 2 \end{array}$$

	$1 +$	$2 +$	$3 +$	$4 +$	$5 +$	
	<b>6</b>	<b>5</b>	<b>4</b>	<b>3</b>	<b>2</b>	

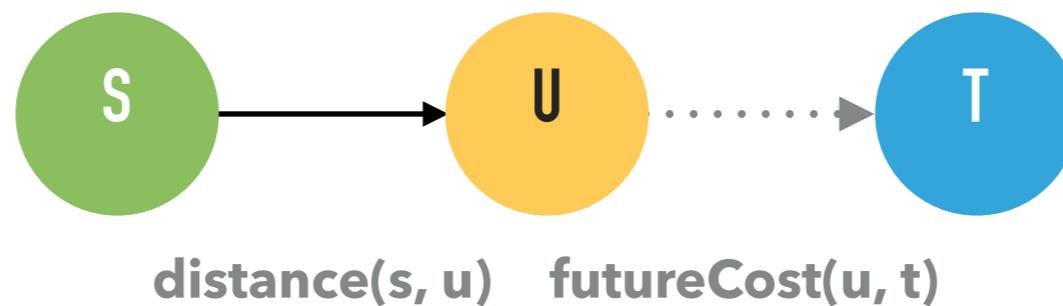
$1 +$	$\star$	1	2	3	4	$\star$
<b>6</b>						

	$1 +$	$2 +$	$3 +$	$4 +$	$5 +$	
	<b>6</b>	<b>5</b>	<b>4</b>	<b>3</b>	<b>2</b>	



**MAKING GOOD  
LIFE DECISIONS**

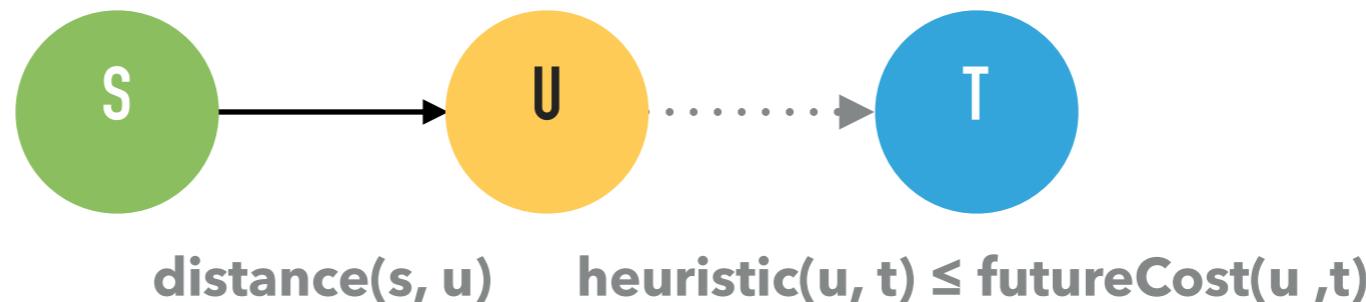
## FORMAL DEFINITIONS



### IDEAL

$priority(u) = distance(s, u)$   
+  $futureCost(u, t)$

## FORMAL DEFINITIONS



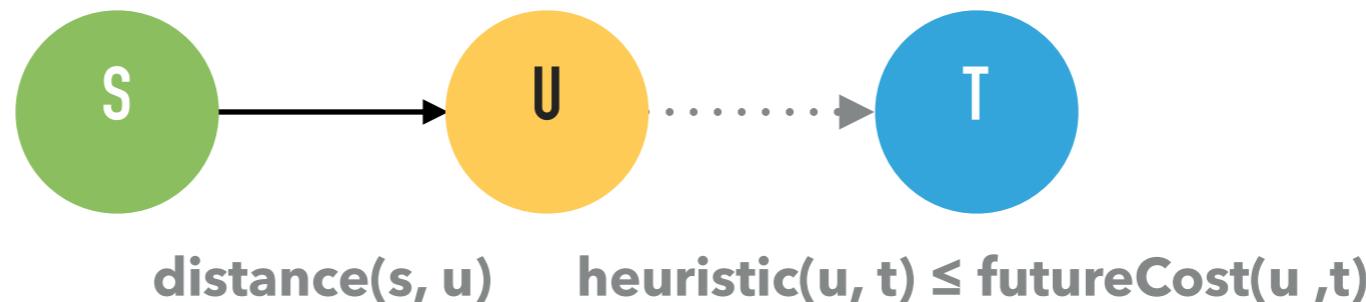
IDEAL

$priority(u) = distance(s, u)$   
+  $futureCost(u, t)$

A\*

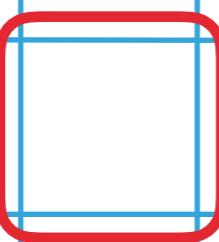
$priority(u) = distance(s, u)$   
+ **heuristic( $u, t$ )**

## HEURISTICS

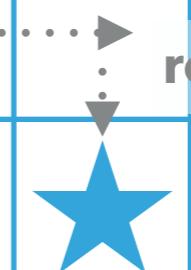


A heuristic is a function that **underestimates** the cost of traveling from  $u$  to  $t$ .

It's a "relaxation" heuristic.



columns apart



rows apart





$1 + 6$

$1 + 6$

$1 + 6$



$1 + 4$



$1 + 6$

$2 + 5$

$1 + 6$

$1 + 6$

$2 + 5$



1

$2 + 3$



$$\begin{array}{c} 1 + \\ 6 \end{array}$$

$$\begin{array}{c} 2 + \\ 5 \end{array}$$

$$\begin{array}{c} 3 + \\ 4 \end{array}$$

$$\begin{array}{c} 1 + \\ 6 \end{array}$$



1

2

$$\begin{array}{c} 1 + \\ 6 \end{array}$$

$$\begin{array}{c} 2 + \\ 5 \end{array}$$

$$\begin{array}{c} 3 + \\ 4 \end{array}$$



**2 +**  
**7**

**2 +**  
**7**

1

**1 +**  
**6**



**2 +**  
**5**

1

**3 +**  
**4**

2



$2 + 7$

$2 + 7$

1

$2 + 5$

$3 + 4$

$2 + 7$

$2 + 7$

1

$2 + 5$

$3 + 4$

$2 + 7$



		$2 + 7$	$3 + 6$			
		$2 + 7$	1	2	$3 + 4$	
	$2 + 7$	1	★	1	2	
		$2 + 7$	1	2	$3 + 4$	
		$2 + 7$	$3 + 6$			★

$$\begin{array}{c} 2+ \\ 7 \end{array}$$

$$\begin{array}{c} 3+ \\ 6 \end{array}$$

$$\begin{array}{c} 4+ \\ 5 \end{array}$$

$$\begin{array}{c} 2+ \\ 7 \end{array}$$

1

2

3

$$\begin{array}{c} 2+ \\ 7 \end{array}$$

1



1

2

$$\begin{array}{c} 2+ \\ 7 \end{array}$$

1

2

3

$$\begin{array}{c} 2+ \\ 7 \end{array}$$

$$\begin{array}{c} 3+ \\ 6 \end{array}$$

$$\begin{array}{c} 4+ \\ 5 \end{array}$$



$2 + 7$

$3 + 6$

$4 + 5$

$3 + 8$

$2 + 7$

1

2

3

$3 + 8$

2

1



1

2



$3 + 8$

$2 + 7$

1

2

3

$2 + 7$

$3 + 6$

$4 + 5$

$$\begin{array}{l} 3+ \\ 8 \end{array}$$
$$\begin{array}{l} 2+ \\ 7 \end{array}$$
$$\begin{array}{l} 3+ \\ 6 \end{array}$$
$$\begin{array}{l} 4+ \\ 5 \end{array}$$

$$\begin{array}{l} 3+ \\ 8 \end{array}$$

2 1 2 3

$$\begin{array}{l} 3+ \\ 8 \end{array}$$

2

1  1 2



$$\begin{array}{l} 3+ \\ 8 \end{array}$$

2 1 2 3

$$\begin{array}{l} 3+ \\ 8 \end{array}$$
$$\begin{array}{l} 2+ \\ 7 \end{array}$$
$$\begin{array}{l} 3+ \\ 6 \end{array}$$
$$\begin{array}{l} 4+ \\ 5 \end{array}$$

$3 + 8$

$3 + 8$

2

$3 + 6$

$4 + 5$

$3 + 8$

2

1

2

3

$3 + 8$

2

1



1

2

2



$3 + 8$

2

1

2

3

$3 + 8$

2

$3 + 6$

$4 + 5$

$3 + 8$

			$3 + 8$	$4 + 7$		
		$3 + 8$	2	3	$4 + 5$	
	$3 + 8$	2	1	2	3	
$3 + 8$	2	1	★	1	2	
$3 + 8$	2	1	2	3		
	$3 + 8$	2	3	$4 + 5$		
		$3 + 8$	$4 + 7$			

			$3 + 8$	$4 + 7$	$5 + 6$	
		$3 + 8$	2	3	4	$5 + 4$
	$3 + 8$	2	1	2	3	
	$3 + 8$	2	1	1	2	
	$3 + 8$	2	1	2	3	
		$3 + 8$	2	3	4	$5 + 4$
		$3 + 8$	$4 + 7$	$5 + 6$		



		$3 + 8$	$4 + 7$	$5 + 6$	$6 + 5$	
	$3 + 8$	2	3	4	5	$6 + 3$
	$3 + 8$	2	1	2	3	
	$3 + 8$	2	1	1	2	
	$3 + 8$	2	1	2	3	
	$3 + 8$	2	3	4	$5 + 4$	
	$3 + 8$	$4 + 7$	$5 + 6$			



$$\begin{array}{cccc} 3+ & 4+ & 5+ & 6+ \\ 8 & 7 & 6 & 5 \end{array}$$

$$\begin{array}{cccc} 3+ & 2 & 3 & 4 \\ 8 & & & \\ & & 5 & 6+ \\ & & & 3 \end{array}$$

$$\begin{array}{cccc} 3+ & 2 & 1 & 2 \\ 8 & & & \\ & & 3 & \\ & & & \end{array}$$

$$\begin{array}{cccc} 3+ & 2 & 1 & \star \\ 8 & & & \\ & & 1 & 2 \\ & & & \end{array}$$

$$\begin{array}{cccc} 3+ & 2 & 1 & 2 \\ 8 & & & \\ & & 3 & \\ & & & \end{array}$$

$$\begin{array}{cccc} 3+ & 2 & 3 & 4 \\ 8 & & & \\ & & 5 & 6+ \\ & & & 3 \end{array}$$

$$\begin{array}{cccc} 3+ & 4+ & 5+ & 6+ \\ 8 & 7 & 6 & 5 \end{array}$$

			$3 + 8$	$4 + 7$	$5 + 6$	$6 + 5$	$6 + 4$		
		$3 + 8$	2	3	4	5	6	$7 + 2$	
	$3 + 8$	2	1	2	3		$7 + 2$		
	$3 + 8$	2	1	★	1	2		★	
	$3 + 8$	2	1	2	3				
		$3 + 8$	2	3	4	5	$6 + 3$		
			$3 + 8$	$4 + 7$	$5 + 6$	$6 + 5$			

		$3 + 8$	$4 + 7$	$5 + 6$	$6 + 5$	$6 + 4$	
		$3 + 8$	2	3	4	5	$7 + 2$
		$3 + 8$	2	1	2	3	$8 + 1$
	$3 + 8$	2	1	★	1	2	$8 + 1$
	$3 + 8$	2	1	2	3		
		$3 + 8$	2	3	4	5	$6 + 3$
		$3 + 8$	$4 + 7$	$5 + 6$	$6 + 5$		

		$3 + 8$	$4 + 7$	$5 + 6$	$6 + 5$	$6 + 4$	
		$3 + 8$	2	3	4	5	$7 + 2$
		$3 + 8$	2	1	2	3	$8 + 1$
	$3 + 8$	2	1	★	1	2	$9 + 0$
	$3 + 8$	2	1	2	3		$9 + 2$
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		$3 + 8$	$4 + 7$	$5 + 6$	$6 + 5$	$6 + 4$	
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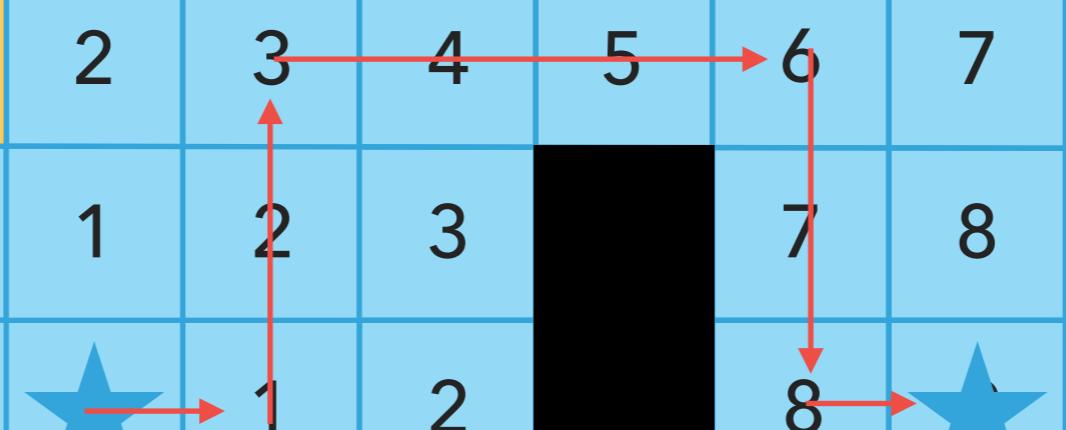
		$3 + 8$	$4 + 7$	$5 + 6$	$6 + 5$	$6 + 4$	$8 + 3$	
		$3 + 8$	2	3	4	5	6	$8 + 3$
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9	9	8	7	6	5	6	7	8	9			
9	8	7	6	5	4	5	6	7	8	9		
8	7	6	5	4	3	4	5	6	7	8	9	
7	6	5	4	3	2	3	4	5	6	7	8	9
6	5	4	3	2	1	2	3		7	8	9	
5	4	3	2	1	★	1	2		8	★		
6	5	4	3	2	1	2	3		7	8	9	
7	6	5	4	3	2	3	4	5	6	9		
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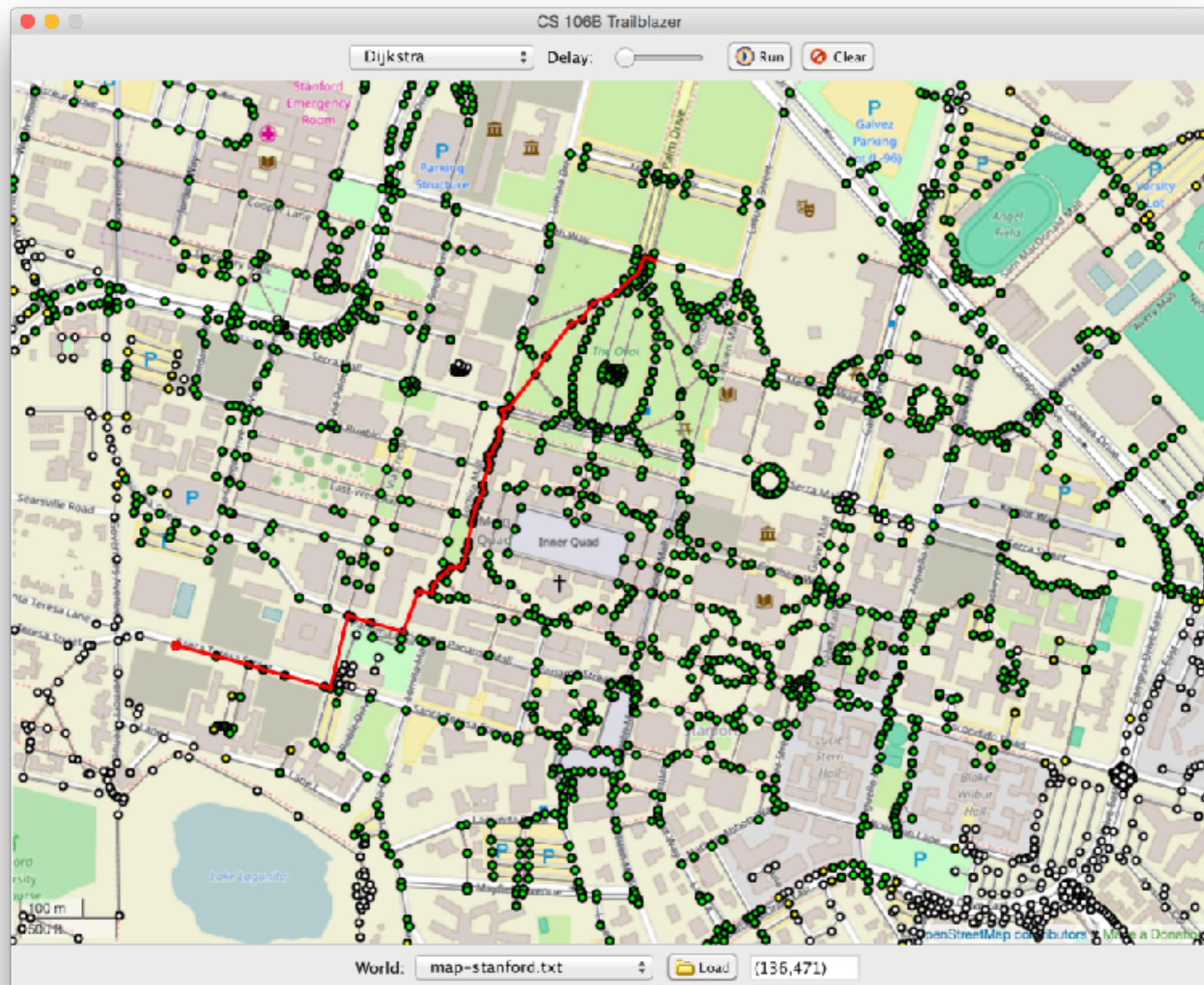
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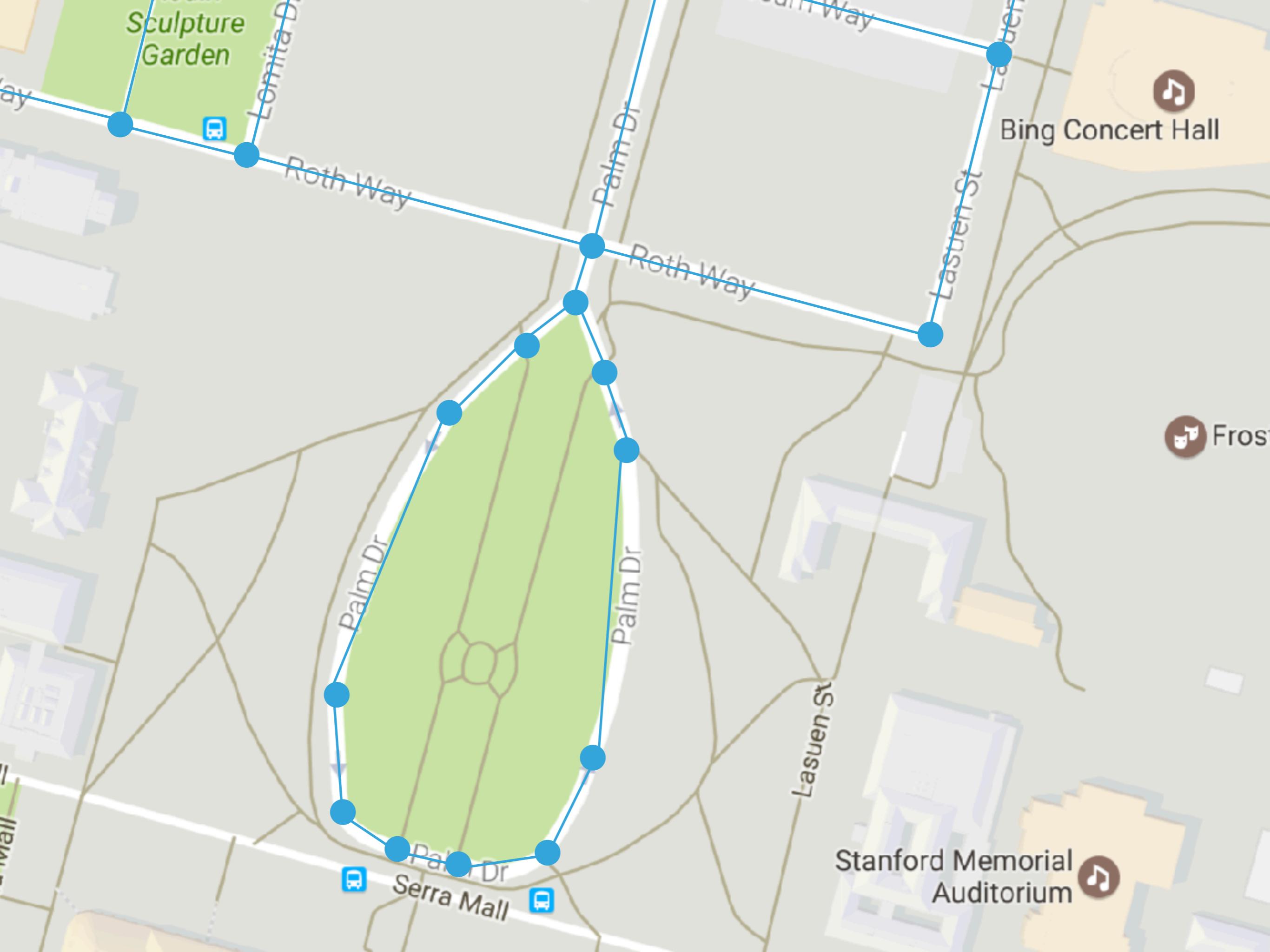
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**YOU WANT YOUR HEURISTIC TO BE AS LARGE AS POSSIBLE**

**BUT YOU NEVER WANT IT TO BE LARGER THAN THE ACTUAL COST.**

# GOOGLE MAPS





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  - ▶ Or... 190 million years

THE LIFE CHANGING MAGIC OF DIJKSTRA AND A\*

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- ▶ All of these and more?
  - ▶ You can use multiple heuristics and choose the max