

CS107 Winter 2020, Lecture 3

Bits and Bytes; Bitwise Operators

reading:

Bryant & O'Hallaron, Ch. 2.1

Announcements

- Assignment 0 deadline tonight at 11:59PM PST
- Assignment 1 (Bit operations!) goes (or rather, went) out today
 - Saturated arithmetic
 - Cell Automata
 - Unicode and UTF-8
- Labs start this week!

Plan For Today

- Bitwise Operators and Masks
- **Demo 1:** Fun With Bits
- **Demo 2:** Powers of 2
- Bit Shift Operators
- **Demo 3:** Color Wheel

Plan For Today

- **Bitwise Operators and Masks**
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**Now that we understand
binary representations, how
can we manipulate them at the
bit level?**

Bitwise Operators

- You're already familiar with many operators in C:
 - **Arithmetic operators:** +, -, *, /, %
 - **Comparison operators:** ==, !=, <, >, <=, >=
 - **Logical Operators:** &&, ||, !
- Today, we're introducing a new category of operators: **bitwise operators:**
 - &, |, ~, ^, <<, >>

And (&)

& is a binary operator. The & of 2 bits is 1 if both bits are 1, and 0 otherwise.

output = a & b;

a	b	output
0	0	0
0	1	0
1	0	0
1	1	1

Or (|)

| is a binary operator. The | of 2 bits is 1 if either (or both) bits is 1.

output = a | b;

a	b	output
0	0	0
0	1	1
1	0	1
1	1	1

Not (\sim)

\sim is a unary operator. The \sim of a bit is 1 if the bit is 0, or 0 otherwise.

output = \sim a;

a	output
0	1
1	0

Exclusive Or (^)

\wedge is a binary operator. The \wedge of 2 bits is 1 iff *exactly* one of the bits is 1.

$$\text{output} = a \wedge b;$$

a	b	output
0	0	0
0	1	1
1	0	1
1	1	0

Operators on Multiple Bits

- When these operators are applied to numbers (multiple bits), the operator is applied to the corresponding bits in each number. For example:


AND	OR	XOR	NOT
<div>0110 & 1100 ---- 0100</div>	<div>0110 1100 ---- 1110</div>	<div>0110 ^ 1100 ---- 1010</div>	<div>~ 1100 ---- 0011</div>

Note: these are different from the logical operators AND (&&), OR (||) and NOT (!).

Operators on Multiple Bits

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This is different from logical AND (&&). The logical AND returns true if both are nonzero, or false otherwise.

Operators on Multiple Bits

- When these operators are applied to numbers (multiple bits), the operator is applied to the corresponding bits in each number. For example:

AND

	0110
&	1100

	0100

OR

	0110
	1100

	1110

XOR


	0110
^	1100

	1010

NOT

~	1100

	0011



This is different from logical NOT (!). The logical NOT returns true if this is zero, and false otherwise.

Bit Vectors and Sets

- We can use bit vectors (ordered collections of bits) to represent finite sets, and perform functions such as union, intersection, and complement.
- **Example:** we can represent current courses taken using a **char**.

0	0	1	0	0	0	1	1
CS161	CS109	CS103	CS110	CS107	CS106X	CS106B	CS106A

Bit Vectors and Sets

0	0	1	0	0	0	1	1
CS161	CS109	CS103	CS110	CS107	CS106X	CS106B	CS106A

- How do we find the union of two sets of courses taken? Use |:

```
    00100011
  | 01100001
  -----
    01100011
```

Bit Vectors and Sets

0	0	1	0	0	0	1	1
CS161	CS109	CS103	CS110	CS107	CS106X	CS106B	CS106A

- How do we find the intersection of two sets of courses taken? Use &:

```
    00100011
&  01100001
-----
    00100001
```


Bit Masking

- We will frequently want to manipulate or isolate out specific bits in a larger collection of bits. A **bitmask** is a constructed bit pattern that we can use, along with bit operators, to do this.
- **Example:** how do we update our bit vector to indicate we've taken CS107?

0	0	1	0	0	0	1	1
CS161	CS109	CS103	CS110	CS107	CS106X	CS106B	CS106A

```
      00100011
    | 00001000
    -----
      00101011
```

Bit Masking

```
#define CS106A 0x1    /* 0000 0001 */
#define CS106B 0x2    /* 0000 0010 */
#define CS106X 0x4    /* 0000 0100 */
#define CS107  0x8    /* 0000 1000 */
#define CS110  0x10   /* 0001 0000 */
#define CS103  0x20   /* 0010 0000 */
#define CS109  0x40   /* 0100 0000 */
#define CS161  0x80   /* 1000 0000 */

char myClasses = ...;
myClasses = myClasses | CS107;    // Add CS107
```

Bit Masking

```
#define CS106A 0x1    /* 0000 0001 */
#define CS106B 0x2    /* 0000 0010 */
#define CS106X 0x4    /* 0000 0100 */
#define CS107  0x8    /* 0000 1000 */
#define CS110  0x10   /* 0001 0000 */
#define CS103  0x20   /* 0010 0000 */
#define CS109  0x40   /* 0100 0000 */
#define CS161  0x80   /* 1000 0000 */
```

```
char myClasses = ...;
myClasses |= CS107;    // Add CS107
```

Bit Masking

- **Example:** how do we update our bit vector to indicate we've *not* taken CS103?

0	0	1	0	0	0	1	1
CS161	CS109	CS103	CS110	CS107	CS106X	CS106B	CS106A

```
    00100011
  & 11011111
  -----
    00000011
```

```
char myClasses = ...;
myClasses = myClasses & ~CS103;  // Remove CS103
```

Bit Masking

- **Example:** how do we update our bit vector to indicate we've *not* taken CS103?

0	0	1	0	0	0	1	1
CS161	CS109	CS103	CS110	CS107	CS106X	CS106B	CS106A

```
    00100011
  & 11011111
  -----
    00000011
```

```
char myClasses = ...;
myClasses &= ~CS103;    // Remove CS103
```

Bit Masking

- **Example:** how do we check if we've taken CS106B?

0	0	1	0	0	0	1	1
CS161	CS109	CS103	CS110	CS107	CS106X	CS106B	CS106A

```
      00100011
    & 00000010
    -----
      00000010
```

```
char myClasses = ...;
if (myClasses & CS106B) {...
    // taken CS106B!
```

Bit Masking

- **Example:** how do we check if we've *not* taken CS107?

0	0	1	0	0	0	1	1
CS161	CS109	CS103	CS110	CS107	CS106X	CS106B	CS106A

```
      00100011
    & 00001000
    -----
    00000000
```

```
char myClasses = ...;
if (!(myClasses & CS107)) {...
    // not taken CS107!
```

Bit Masking

- **Example:** how do we check if we've *not* taken CS107?

0	0	1	0	0	0	1	1
CS161	CS109	CS103	CS110	CS107	CS106X	CS106B	CS106A

	00100011		00000000
&	00001000	^	00001000
	-----		-----
	00000000		00001000

```
char myClasses = ...;  
if ((myClasses & CS107) ^ CS107) {...  
    // not taken CS107!
```


Demo: Fun with Bits in gdb



Bit Masking

Bit masking is also useful for integer representations as well. For instance, we might want to check the value of the most-significant bit, or just one of the middle bytes.

- **Example:** If I have a 32-bit integer **j**, what operation should I perform if I want to extract *just the lowest byte* in **j**?

```
int j = ...;  
int k = j & 0xff;           // mask to get just lowest byte
```

Practice: Bit Masking

- **Practice 1:** write an expression that, given a 32-bit integer j , sets its least-significant byte to all 1s, but preserves all other bytes.

$j \mid 0xff$

- **Practice 2:** write an expression that, given a 32-bit integer j , flips ("complements") all but the least-significant byte, and preserves all other bytes.

$j \wedge \sim 0xff$

Powers of 2

Without using loops, how can we detect if a binary number is a power of 2? What is special about its binary representation and how can we leverage that?

Demo: Powers of 2



Plan For Today

- Bitwise Operators and Masks
- **Demo 1:** Courses
- **Demo 2:** Powers of 2
- **Bit Shift Operators**

Left Shift (<<)

The LEFT SHIFT operator shifts a bit pattern a certain number of positions to the left. New lower order bits are filled in with 0s, and bits shifted off the end are lost.

```
x << k;    // shifts x to the left by k bits
```

8-bit examples:

```
00110111 << 2 results in 11011100
```

```
01100011 << 4 results in 00110000
```

```
10010101 << 4 results in 01010000
```

Right Shift (>>)

The RIGHT SHIFT operator shifts a bit pattern a certain number of positions to the right. Bits shifted off the end are lost.

```
x >> k;    // shifts x to the right by k bits
```

Question: how should we fill in new higher-order bits?

Idea: let's follow left-shift and fill with 0s.

```
short x = 2;           // 0000 0000 0000 0010
int y = x >> 1;         // 0000 0000 0000 0001
printf("%d\n", y);     // 1
```


Right Shift (>>)

The RIGHT SHIFT operator shifts a bit pattern a certain number of positions to the right. Bits shifted off the end are lost.

```
x >> k;    // shifts x to the right by k bits
```

Question: how should we fill in new higher-order bits?

Idea: let's follow left-shift and fill with 0s.

```
short x = -2;           // 1111 1111 1111 1110
int y = x >> 1;          // 0111 1111 1111 1111
printf("%d\n", y);      // 32767!
```

Right Shift (>>)

The RIGHT SHIFT operator shifts a bit pattern a certain number of positions to the right. Bits shifted off the end are lost.

```
x >> k;    // shifts x to the right by k bits
```

Question: how should we fill in new higher-order bits?

Problem: always filling with zeros means we may change the sign bit.

Solution: let's fill with the sign bit!

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```
x >> k;    // shifts x to the right by k bits
```

Question: how should we fill in new higher-order bits?

Solution: let's fill with the sign bit!

```
short x = 2;           // 0000 0000 0000 0010
int y = x >> 1;        // 0000 0000 0000 0001
printf("%d\n", y);     // 1
```

Right Shift (>>)

The RIGHT SHIFT operator shifts a bit pattern a certain number of positions to the right. Bits shifted off of the end are lost.

```
x >> k;    // shifts x to the right by k bits
```

Question: how should we fill in new higher-order bits?

Solution: let's fill with the sign bit!

```
short x = -2;           // 1111 1111 1111 1110
int y = x >> 1;          // 1111 1111 1111 1111
printf("%d\n", y);      // -1!
```

Right Shift (>>)

There are *two kinds* of right shifts, depending on the value and type you are shifting:

- **Logical Right Shift:** fill new high-order bits with 0s.
- **Arithmetic Right Shift:** fill new high-order bits with the most-significant bit.

Unsigned numbers are right-shifted using **Logical Right Shift**.

Signed numbers are right-shifted using **Arithmetic Right Shift**.

This way, the sign of the number (if applicable) is preserved!

Shift Operator Pitfalls

1. *Technically*, the C standard does not precisely define whether a right shift for signed integers is logical or arithmetic. However, **almost all compilers/machines** use arithmetic, and you can most likely assume this.
2. Operator precedence can be tricky! For example:

$1 \ll 2 + 3 \ll 4$ means $1 \ll (2+3) \ll 4$ because addition and subtraction have higher precedence than shifts! Always use parentheses to be sure:

$(1 \ll 2) + (3 \ll 4)$

Shift Operator Pitfalls

- The default type of a number literal in your code is an **int**.
- Let's say you want a long with the index-32 bit as 1:

```
long num = 1 << 32;
```

- This doesn't work! 1 is by default an **int**, and you can't shift an int by 32 because it only has 32 bits. You must specify that you want 1 to be a **long**.

```
long num = 1L << 32;
```

Demo: Color Wheel



Recap

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Next time: *How can a computer represent and manipulate more complex data like text?*

Bonus Demo: Bitwise abs

