

CS107 Lecture 4

Bits and Bytes; Bitwise Operators

reading:

Bryant & O'Hallaron, Ch. 2.1



masks recommended

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Based on slides created by Cynthia Lee, Chris Gregg, Jerry Cain, Lisa Yan and others.

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CS107 Topic 1

How can a computer represent integer numbers?

Why is answering this question important?

- Helps us understand the limitations of computer arithmetic (last time)
- Shows us how to more efficiently perform arithmetic (today)
- Shows us how we can encode data more compactly and efficiently (today)

assign1: implement 3 programs that manipulate binary representations to (1) work around the limitations of arithmetic with addition, (2) simulate an evolving colony of cells, and (3) print Unicode text to the terminal.

Today

Today, we'll learn about a new set of operators to manipulate bits. For example:

```
int x = 2;
```

```
// NEW: shift all bits x places to the left or right
x = x << 1; // now x is 4!
```

```
// NEW: check if the least significant bit is a 0
if (x & 1 == 0) {...
```

This is useful because we can perform some arithmetic more efficiently, and also store data more compactly in individual bits.

Learning Goals

- Learn about the bitwise C operators and how to use them to manipulate bits
- Understand when to use one bitwise operator vs. another in your program
- Get practice with writing programs that manipulate binary representations

Lecture Plan

- **Recap and continuing:** Integer Representations
- Bitwise Operators
- Bitmasks

Lecture Plan

- **Recap and continuing: Integer Representations**
 - Bitwise Operators
 - Bitmasks

Bits and Bytes So Far

All data, including integer variables, are ultimately stored in memory in binary:

```
int x = 5;      // really 0b0...0101 in memory!
```

- Unsigned numbers store the direct binary representation of its value
- Signed numbers use **two's complement** to store its positive/negative/0 value
- Overflow occurs when we exceed the the minimum or maximum value of the bit representation – it can cause some funky bugs!

Base 10 vs. Binary vs. Hex

- Let's take a byte (8 bits):

165

Base-10: Human-readable,
but cannot easily interpret on/off bits

0b10100101

Base-2: Yes, computers use this,
but not human-readable

0xa5

Base-16: Easy to convert to Base-2,
More “portable” as a human-readable format
(fun fact: a half-byte is called a nibble or nybble)

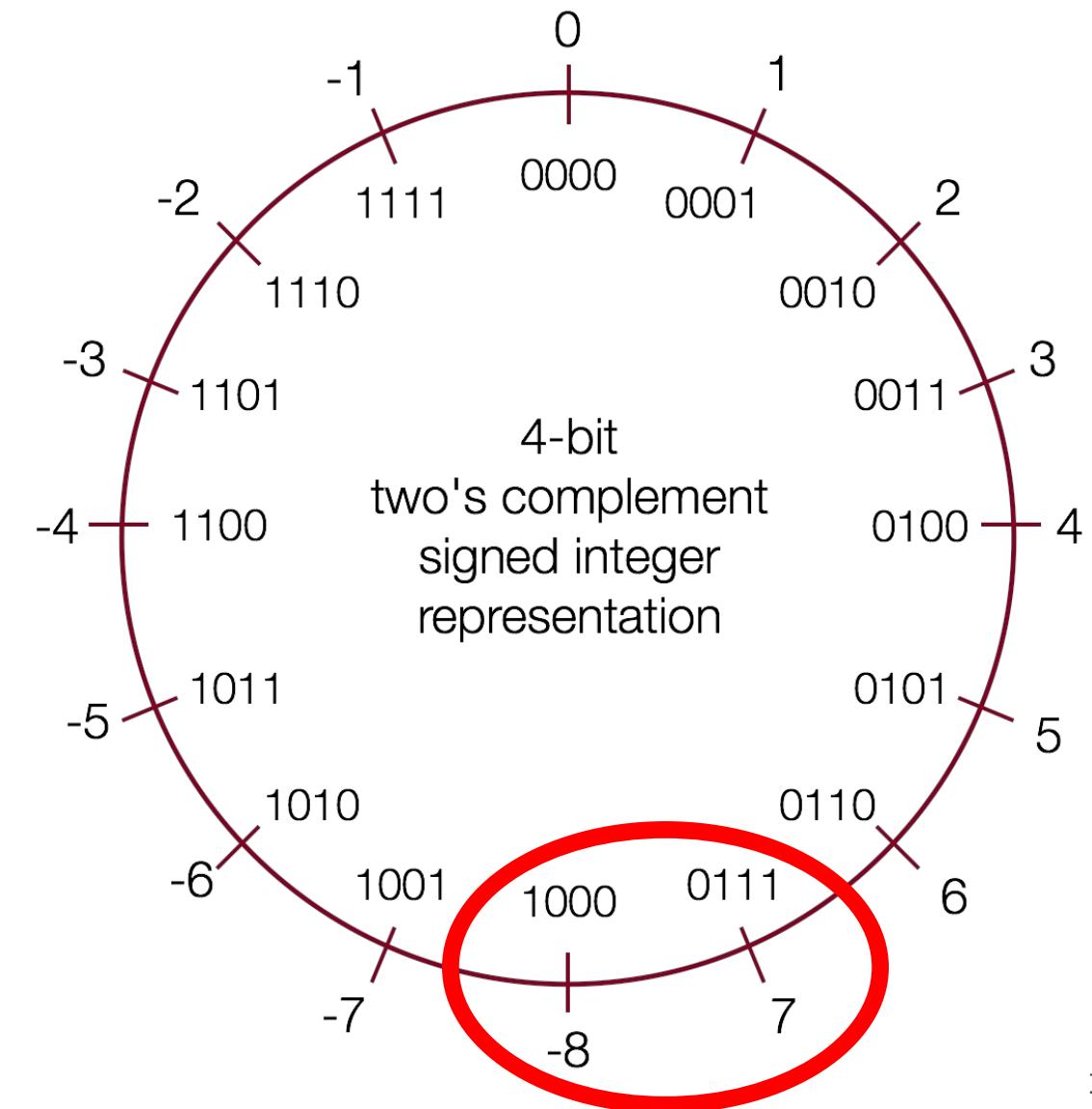
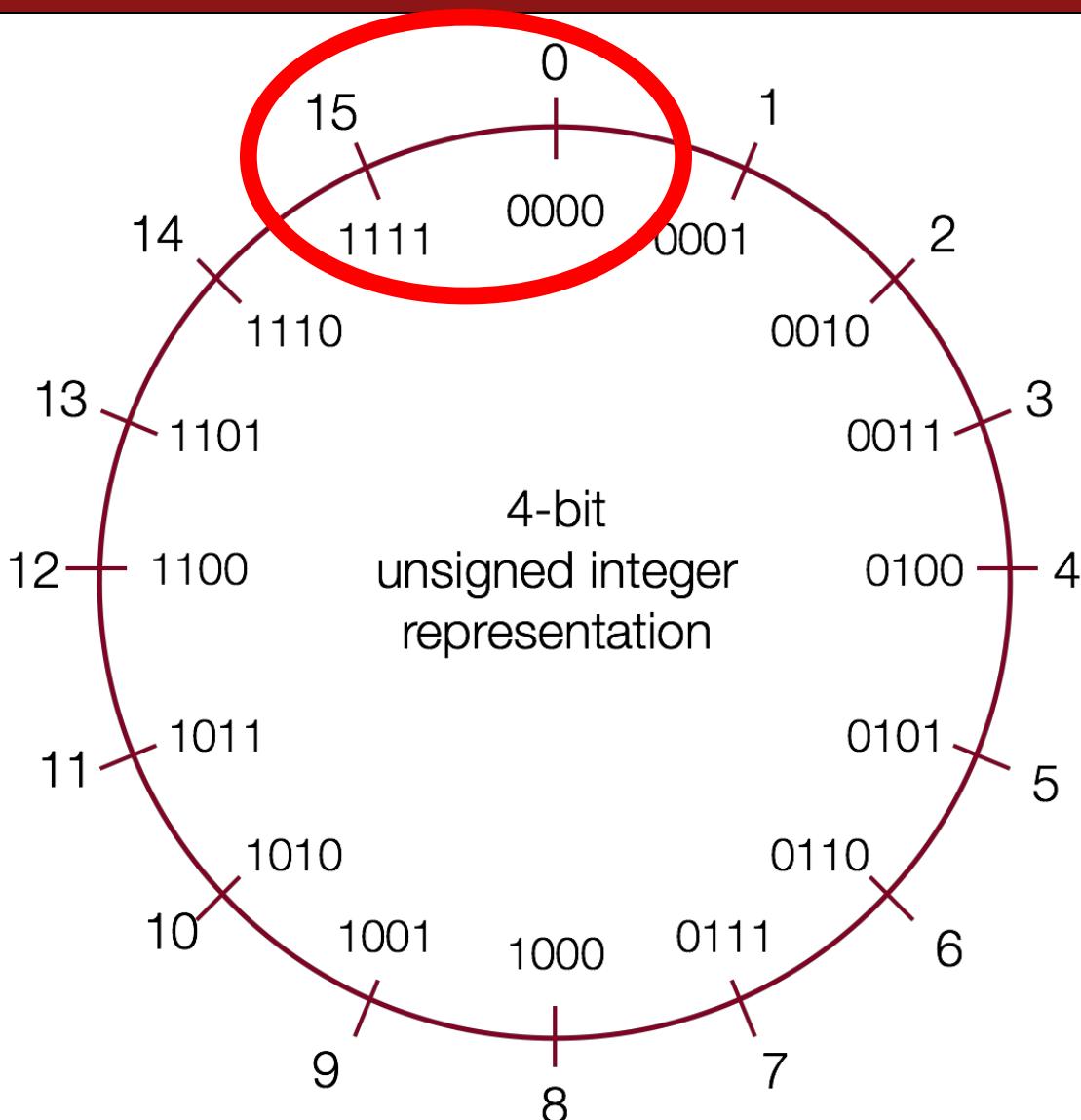
Overflow

Overflow occurs because we don't have enough bits to store a value.

E.g. if we have **unsigned short x = 65535** and add 2, we get **1!**

$$\begin{array}{r} 1111 & 1111 & 1111 & 1111 \\ 0000 & 0000 & 0000 & 0010 \\ + & & & \\ \hline \cancel{1}0000 & 0000 & 0000 & 0001 \end{array}$$

Overflow



Min and Max Integer Values

In C, there are various constants that represent these minimum and maximum values: `INT_MIN`, `INT_MAX`, `UINT_MAX`, `LONG_MIN`, `LONG_MAX`, `ULONG_MAX`, ...

Truncating Bit Representation

If we want to **reduce** the bit size of a number, C *truncates* the representation and discards the *more significant bits*.

```
int x = 53191;  
short sx = x; // -12345!
```

```
x = 0000 0000 0000 0000 1100 1111 1100 0111  
sx = 1100 1111 1100 0111
```

Truncating Bit Representation

If we want to **reduce** the bit size of a number, C *truncates* the representation and discards the *more significant bits*.

```
int x = -3;  
short sx = x; // still -3
```

```
x = 1111 1111 1111 1111 1111 1111 1111 1101  
sx = 1111 1111 1111 1101
```

Expanding Bit Representations

Sometimes, we want to carry over a value to a larger variable (e.g. make an **int** and set it equal to a **short**).

- For **unsigned** values, C adds *leading zeros* to the representation (“zero extension”)
- For **signed** values, C *repeats the sign of the value* for new digits (“sign extension”)

Expanding Bit Representation

If we want to **expand** the bit size of an **unsigned** number, C *adds leading zeros*.

```
unsigned short s = 4;  
unsigned int i = s; // still 4
```

s =	0000 0000 0000 0100
i =	0000 0000 0000 0000 0000 0000 0100

Expanding Bit Representation

If we want to **expand** the bit size of an **signed** number, C *adds repeats the sign*.

```
short s = -4;  
int i = s; // still -4
```

s =	1111 1111 1111 1100
i =	1111 1111 1111 1111 1111 1111 1111 1100

Expanding Bit Representation

If we want to **expand** the bit size of an **signed** number, C *adds repeats the sign*.

```
short s = 4;  
int i = s; // still 4
```

s =	0000 0000 0000 0100
i =	0000 0000 0000 0000 0000 0000 0100

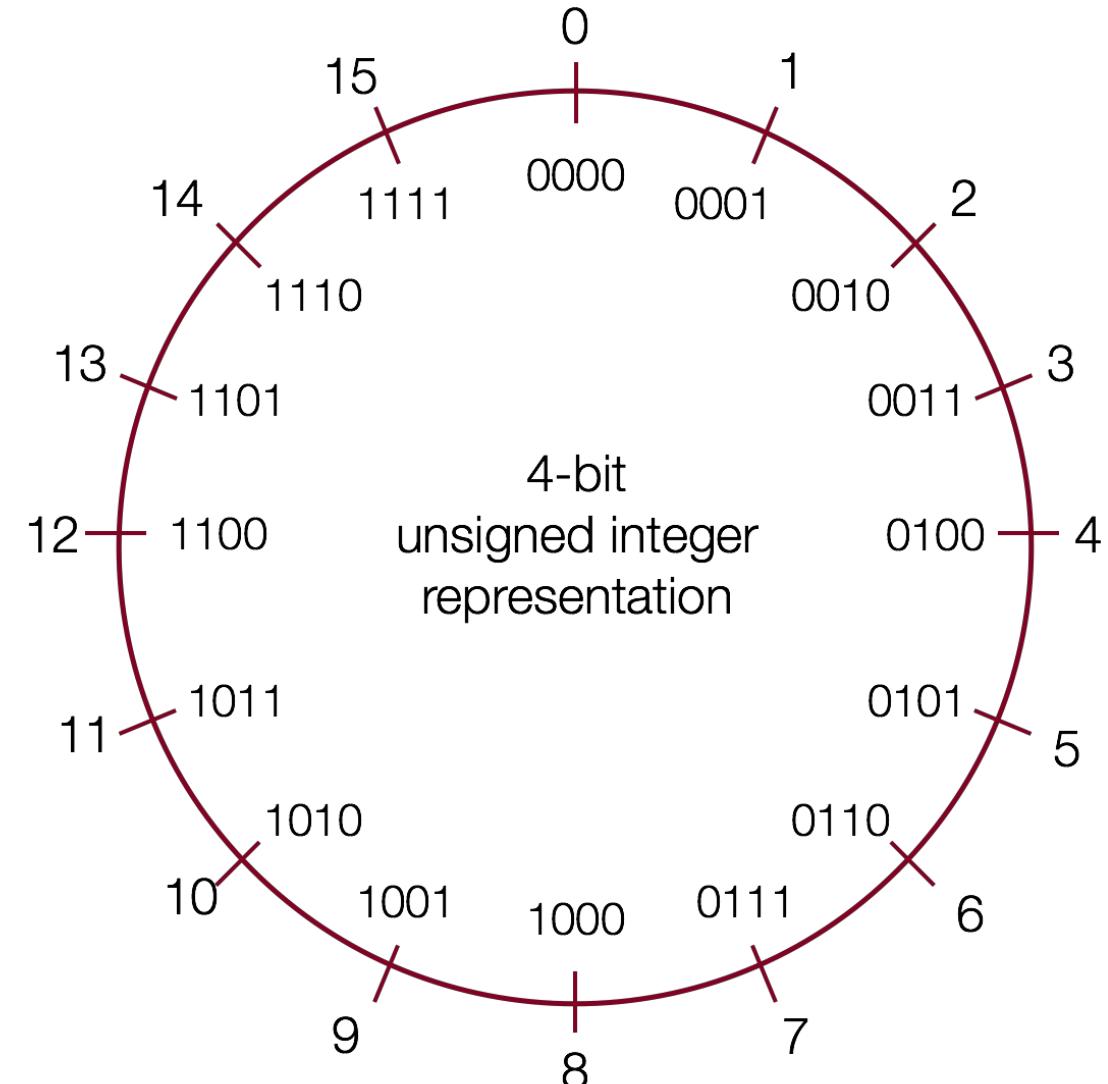
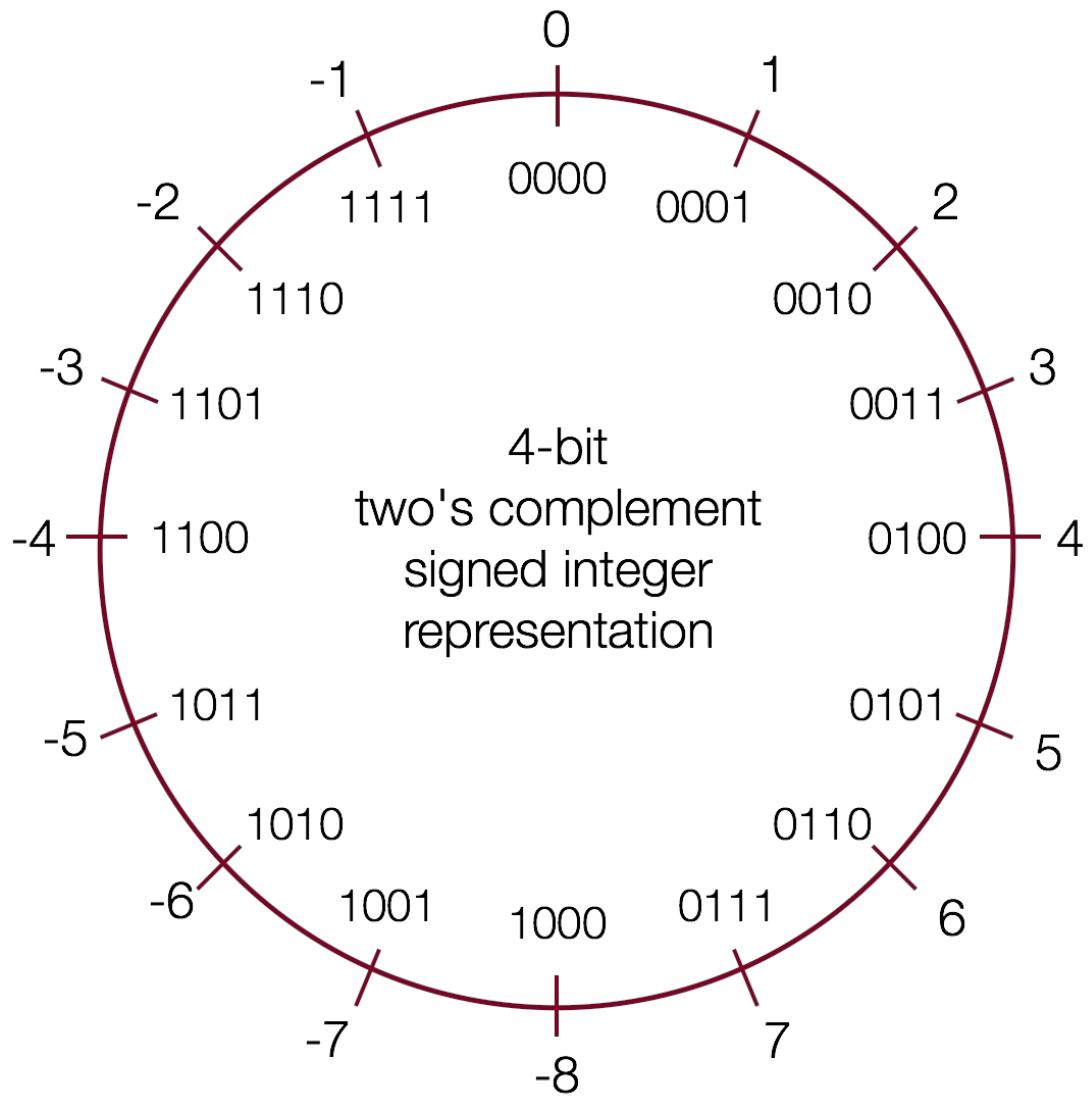
Casting

Casting between variable types can cause tricky issues; **the bits remain the same but are interpreted differently.**

```
short s = -12345;  
unsigned short us = (unsigned short)s; // 53191!
```

Both **s** and **us** have the same binary representation: **1100 1111 1100 0111**
But its represented value varies depending on how it's interpreted!

Casting



Casting

You can cast something to another type by putting that type in parentheses in front of the value:

```
int v = -12345;  
... (unsigned int)v...
```

You can also use the **U** suffix after a number literal to treat it as unsigned:

-12345U

Comparisons Between Different Types

Be **careful** when comparing signed and unsigned integers. **C will implicitly cast** the signed argument to unsigned, and then performs the operation assuming both numbers are non-negative.

```
int x = -1;      // 1111 1111 1111 1111 1111 1111 1111 1111 1111
unsigned int y = 0;
if (x < y) { ... // will be false!!
```

Note: when doing `<`, `>`, `<=`, `>=` comparison between different size types, it will *promote to the larger type*.

The sizeof Operator

```
long sizeof(type);
```

// Example

```
long int_size_bytes = sizeof(int);    // 4
long short_size_bytes = sizeof(short); // 2
long char_size_bytes = sizeof(char);   // 1
```

sizeof takes a variable type (or a variable itself) as a parameter and returns the size of that type, in bytes.

Lecture Plan

- Recap and continuing: Integer Representations
- **Bitwise Operators**
- Bitmasks

Bitwise Operators

- You're already familiar with many operators in C:
 - **Arithmetic operators:** +, -, *, /, %
 - **Comparison operators:** ==, !=, <, >, <=, >=
 - **Logical Operators:** &&, ||, !
- Today, we're introducing a new category of operators: **bitwise operators:**
 - &, |, ~, ^, <<, >>

And (&)

AND is a binary operator. The AND of 2 bits is 1 if both bits are 1, and 0 otherwise.

output = a & b;

a	b	output
0	0	0
0	1	0
1	0	0
1	1	1

& with 1 to let a bit through, & with 0 to zero out a bit

Or (|)

OR is a binary operator. The OR of 2 bits is 1 if either (or both) bits is 1.

output = a | b;

a	b	output
0	0	0
0	1	1
1	0	1
1	1	1

| with 1 to turn on a bit, | with 0 to let a bit go through

Not (\sim)

NOT is a unary operator. The NOT of a bit is 1 if the bit is 0, or 1 otherwise.

output = $\sim a$;

a	output
0	1
1	0

Exclusive Or (^)

Exclusive Or (XOR) is a binary operator. The XOR of 2 bits is 1 if *exactly* one of the bits is 1, or 0 otherwise.

output = a ^ b;

a	b	output
0	0	0
0	1	1
1	0	1
1	1	0

\wedge with 1 to flip a bit, \wedge with 0 to let a bit go through

Operators on Multiple Bits

When these operators are applied to numbers (multiple bits), the operator is applied to the corresponding bits in each number. For example:

AND

$$\begin{array}{r} 0110 \\ \& 1100 \\ \hline 0100 \end{array}$$

OR

$$\begin{array}{r} 0110 \\ | 1100 \\ \hline 1110 \end{array}$$

XOR

$$\begin{array}{r} 0110 \\ ^ 1100 \\ \hline 1010 \end{array}$$

NOT

$$\begin{array}{r} \sim 1100 \\ \hline 0011 \end{array}$$

Bit Operators

```
int x = 4;    // 0000 ... 0100
int y = 5;    // 0000 ... 0101

// 4
int anded = x & y; // 0000 ... 0100

// 5
int ored  = x | y; // 0000 ... 0101

// -5
int notX  = ~x; // 1111 1111 1111 1011

int xored = x ^ y; // what would this give us?
```

When poll is active, respond at **pollev.com/cs107**

 Text **CS107** to **22333** once to join

If $x = 4$ and $y = 5$, what would x^y be?

4

5

1

0

2

If $x = 4$ and $y = 5$, what would x^y be?

4
5
1
0
2

If $x = 4$ and $y = 5$, what would x^y be?

4
5
1
0
2

Operators on Multiple Bits

When these operators are applied to numbers (multiple bits), the operator is applied to the corresponding bits in each number. For example:

AND

$$\begin{array}{r} 0110 \\ \& 1100 \\ \hline 0100 \end{array}$$

OR

$$\begin{array}{r} 0110 \\ | 1100 \\ \hline 1110 \end{array}$$

XOR

$$\begin{array}{r} 0110 \\ ^ 1100 \\ \hline 1010 \end{array}$$

NOT

$$\begin{array}{r} \sim 1100 \\ \hline 0011 \end{array}$$

Note: these are different from the logical operators AND (&&), OR (||) and NOT (!).

Bit Operators

```
int x = 4; // 0000 ... 0100
int y = 5; // 0000 ... 0101
```

```
// This is checking if x and y are both nonzero
if (x && y) { ... }
```

```
// This is checking if the result of x & y is nonzero
if (x & y) { ... }
```

Lecture Plan

- Recap and continuing: Integer Representations
- Bitwise Operators
- **Bitmasks**
- Demo: Bitmasks and GDB
- More practice

Bitmasks

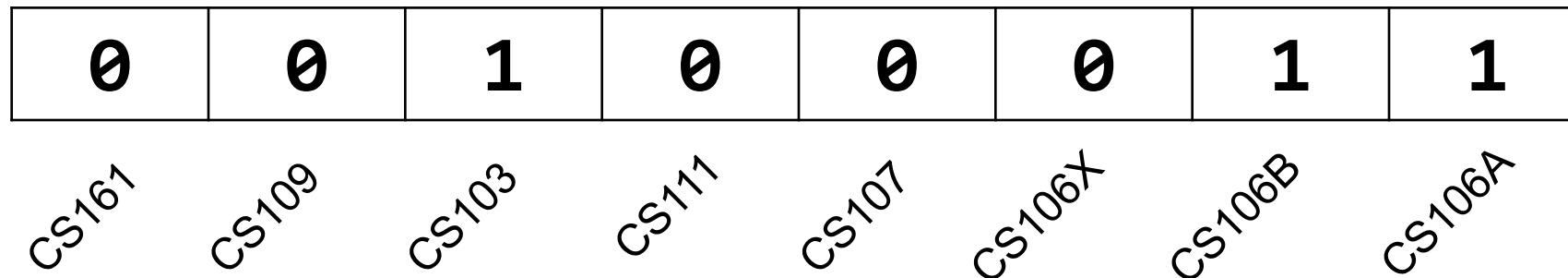
We will frequently want to manipulate or isolate out specific bits in a larger collection of bits. A **bitmask** is a constructed bit pattern that we can use, along with bit operators, to do this.

Motivating Example: Bit vectors

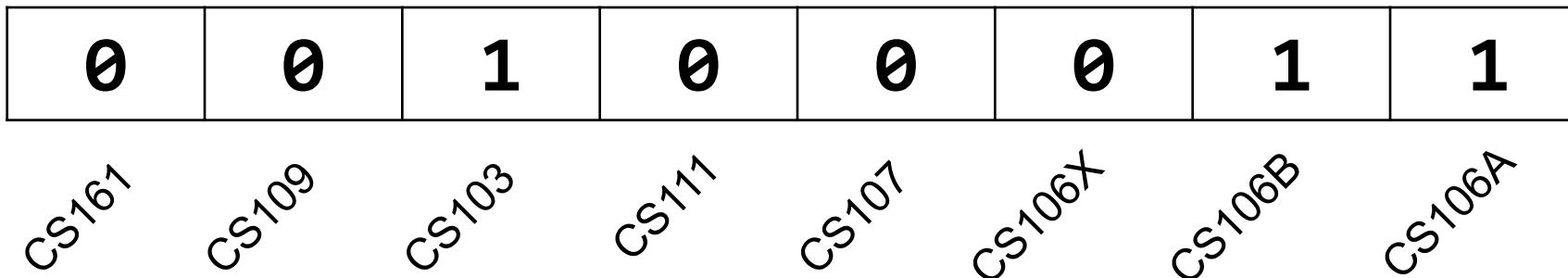
Bit Vectors and Sets

Instead of using arrays of e.g., Booleans in our programs, sometimes it's beneficial to store that information in bits instead – more compact.

- **Example:** we can represent current courses taken using a **char** and manipulate its contents using bit operators.



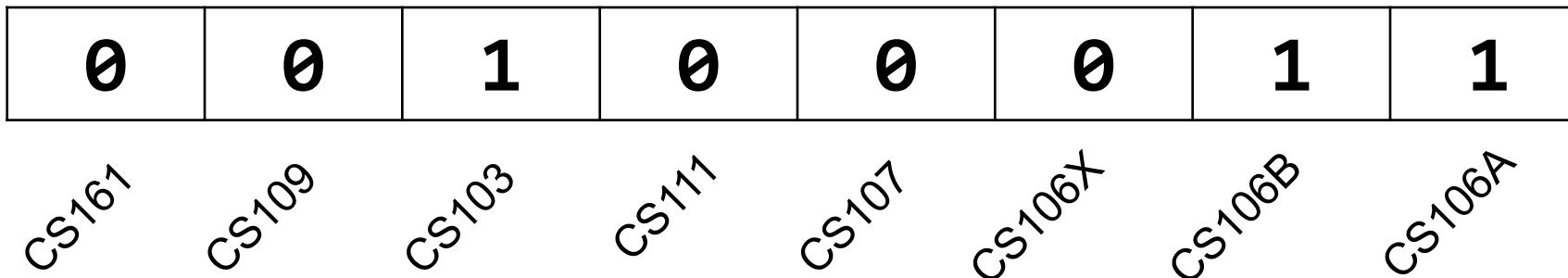
Bit Vectors and Sets



- How do we find the union of two sets of courses taken? Use OR:

$$\begin{array}{r} 00100011 \\ | \quad 01100001 \\ \hline 01100011 \end{array}$$

Bit Vectors and Sets



- How do we find the intersection of two sets of courses taken? Use AND:

$$\begin{array}{r} 00100011 \\ \& 01100001 \\ \hline 00100001 \end{array}$$

Bit Masking

Example: how do we update our bit vector to indicate we've taken CS107?

0	0	1	0	0	0	1	1
CS161	CS109	CS103	CS111	CS107	CS106X	CS106B	CS106A
00100011							
	00001000						

			00101011				

Bit Masking

```
#define CS106A 0x1      /* 0000 0001 */
#define CS106B 0x2      /* 0000 0010 */
#define CS106X 0x4      /* 0000 0100 */
#define CS107  0x8      /* 0000 1000 */
#define CS111 0x10     /* 0001 0000 */
#define CS103 0x20     /* 0010 0000 */
#define CS109 0x40     /* 0100 0000 */
#define CS161 0x80     /* 1000 0000 */

char myClasses = ...;
myClasses = myClasses | CS107;    // Add CS107
```

Bit Masking

```
#define CS106A 0x1      /* 0000 0001 */
#define CS106B 0x2      /* 0000 0010 */
#define CS106X 0x4      /* 0000 0100 */
#define CS107  0x8      /* 0000 1000 */
#define CS111 0x10     /* 0001 0000 */
#define CS103 0x20     /* 0010 0000 */
#define CS109 0x40     /* 0100 0000 */
#define CS161 0x80     /* 1000 0000 */

char myClasses = ...;
myClasses |= CS107;      // Add CS107
```

Bit Masking

- **Example:** how do we update our bit vector to indicate we've *not* taken CS103?

0	0	1	0	0	0	1	1
CS161	CS109	CS103	CS111	CS107	CS106X	CS106B	CS106A

00100011
& 11011111

00000011

```
char myClasses = ...;  
myClasses = myClasses & ~CS103; // Remove CS103
```

Bit Masking

- **Example:** how do we update our bit vector to indicate we've *not* taken CS103?

0	0	1	0	0	0	1	1
CS161	CS109	CS103	CS111	CS107	CS106X	CS106B	CS106A

00100011
& 11011111

00000011

```
char myClasses = ...;  
myClasses &= ~CS103; // Remove CS103
```

Bit Masking

- **Example:** how do we check if we've taken CS106B?

0	0	1	0	0	0	1	1
CS161	CS109	CS103	CS111	CS107	CS106X	CS106B	CS106A

00100011
& 00000010

00000010

```
char myClasses = ...;  
if (myClasses & CS106B) {...  
    // taken CS106B!
```

Bit Masking

- **Example:** how do we check if we've *not* taken CS107?

0	0	1	0	0	0	1	1
CS161	CS109	CS103	CS111	CS107	CS106X	CS106B	CS106A

00100011
& 00001000

00000000

```
char myClasses = ...;  
if (!(myClasses & CS107)) {...  
    // not taken CS107!
```

Bitwise Operator Tricks

- `|` with 1 is useful for turning select bits on
- `&` with 0 is useful for turning select bits off
- `|` is useful for taking the union of bits
- `&` is useful for taking the intersection of bits
- `^` is useful for flipping select bits
- `~` is useful for flipping all bits

Introducing GDB

Is there a way to step through the execution of a program and print out its values as it's running? E.g., to view binary representations? **Yes!**

The GDB Debugger

GDB is a **command-line debugger**, a text-based debugger with similar functionality to other debuggers you may have used, such as in Qt Creator

- It lets you put **breakpoints** at specific places in your program to pause there
- It lets you step through execution line by line
- It lets you print out values of variables in various ways (including binary)
- It lets you track down where your program crashed
- And much, much more!

GDB is essential to your success in CS107 this quarter! We'll be building our familiarity with GDB over the course of the quarter.

Recap

- **Recap and continuing:** Integer Representations
- Bitwise Operators
- Bitmasks

Lecture 3 takeaways: We can use bit operators like `&`, `|`, `~`, etc. to manipulate the binary representation of values. A number is a bit pattern that can be manipulated arithmetically or bitwise at your convenience!

Extra Practice

Color Wheel

- Another application for storing data efficiently in binary is representing **colors**.
- A color representation commonly consists of opacity (how transparent or opaque it is), and how much red/green/blue is in the color.
- **Key idea:** we can encode each of these in 1 byte, in a value from 0-255! Thus, an entire color can be represented in one 4-byte **integer**.

0x 42 53 01 44



Opacity Red Green Blue

Demo: Color Wheel



color_wheel.c

Hexadecimal and Truncation

For each initialization of x, what will be printed?

i. `x = 130; // 0x82`

ii. `x = -132; // 0xff7c`

iii. `x = 25; // 0x19`

```
short x = ____;  
char cx = x;  
printf("%d", cx);
```



Hexadecimal and Truncation

For each initialization of x, what will be printed?

-126 i. `x = 130; // 0x82`

124 ii. `x = -132; // 0xff7c`

```
short x = ____;  
char cx = x;  
printf("%d", cx);
```

25 iii. `x = 25; // 0x19`

Limits and Comparisons

2. Will the following char comparisons evaluate to true or false?

i. $-7 < 4$ **true** iii. $(\text{char}) 130 > 4$ **false**

ii. $-7 < 4U$ **false** iv. $(\text{char}) -132 > 2$ **true**

By default, numeric constants in C are signed ints, unless they are suffixed with u (unsigned) or L (long).

Bitwise Warmup

How can we use bitmasks + bitwise operators to...

0b00001101

1. ...turn **on** a particular set of bits?

0b00001101

0b00001111

2. ...turn **off** a particular set of bits?

0b00001101

0b00001001

3. ...**flip** a particular set of bits?

0b00001101

0b00001011



Bitwise Warmup

How can we use bitmasks + bitwise operators to...

0b00001101

1. ...turn **on** a particular set of bits? **OR**

0b00001101

0b00000010 |

0b00001111

2. ...turn **off** a particular set of bits? **AND**

0b00001101

0b11111011 &

0b00001001

3. ...**flip** a particular set of bits? **XOR**

0b00001101

0b00000110 ^

0b00001011