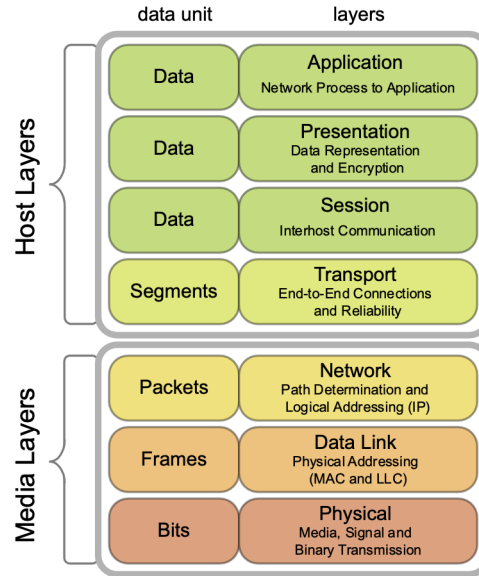


# CS107, Lecture 20: Sockets Programming

# Quick Overview



CultureDuQ, CC BY-SA 4.0 <https://creativecommons.org/licenses/by-sa/4.0>, via

Wikimedia Commons

# TCP and UDP

- Both run on top of IP
- Both have a port number (16 bits)
  - Official port usage is assigned by IANA
  - Ports under 1024 are typically reserved (i.e., on the myth machines, you need special permission to bind to them)
  - Common ports include: 22 (SSH), 53 (DNS), 80 (HTTP), 443 (HTTPS)—see also <https://www.iana.org/assignments/service-names-port-numbers/service-names-port-numbers.xhtml> or `/etc/services`

# TCP and UDP (continued)

- So, to connect to a remote server, we need both an IP address and a port number
  - Quick aside on IP addresses: IPv4 addresses are only 32 bits long, which only makes for about 4 billion total IPs, which we've fully allocated; IPv6 addresses, by contrast, are 128 bits long
    - IPv4 addresses are often written in dotted quad format of 192.168.1.1
    - IPv6 addresses are a little more complicated, but can be written as 2607:f6d0:0:0:0:0:0:0 (which can also be written as 2607:f6d0::)
    - Either way, IP addresses can be a little cumbersome to write, so we use DNS (domain name service) to map from domain names (e.g., web.stanford.edu) to IP addresses (e.g., 171.67.215.200)

# TCP and UDP (continued)

- Quick aside on client server model: for the rest of the lecture, we'll be implicitly referencing this model—the rough idea is that we have a server (imagine, e.g., Google) which serves data to one or more clients (imagine, e.g., people Googling)
- TCP and UDP both allow us to send arbitrary bytes over the network
- It is important that we send bytes in a way that both the client and server will understand
  - A protocol specifies how the bytes will be interpreted
  - IP and TCP/UDP level details specify that the network byte order should be big-endian (myth machines are little-endian)

# TCP and UDP (continued)

- TCP provides a “reliable bytestream” abstraction (except in exceptional cases, the data will arrive correctly on the other side)
  - Useful for non-time critical applications (e.g., web servers (HTTP prior to HTTP/3 runs over TCP), ssh, etc...)
- UDP provides an unreliable datagram abstraction (it’s effectively just a userspace wrapper around IP, hence “User Datagram Protocol”)
  - Useful for time critical applications, or applications that can deal with some data loss (e.g., video conferencing, online gaming, etc...)

# **SOCKET PROGRAMMING BASICS**

# socket ( )

- `int socket(int domain, int type, int protocol);`
- The domain specifies what type of socket we want—for this lecture, it will be one of `PF_INET` or `PF_INET6`
- The type for this lecture will always be `SOCK_STREAM` (meaning TCP, it could also be `SOCK_DGRAM` for UDP)
- The protocol is the protocol number (e.g., one of `IPPROTO_TCP` or `IPPROTO_UDP`, but we can use 0 since `SOCK_STREAM` means TCP, and it will figure it out)
- Returns a “file descriptor” on success and `<0` on error (setting `errno` as appropriate)



# Detour: file descriptors

- You may encounter the phrase “everything is a file” when working in a Unix/Linux context
- File descriptors are one incarnation of this—a `FILE *` is a convenient wrapper around a file descriptor
  - A file descriptor is an integer that the OS hands to our process that we can use syscalls on to read/write data (e.g., `read`, `write`) or otherwise modify (e.g., `fcntl`)
  - We’ll have the following file descriptors always by default: 0 (`stdin`), 1 (`stdout`), 2 (`stderr`)
- Note that we use file descriptors for both real files and for sockets (among other things)

# Detour: error handling

- At Adam's request
- Many system calls (and wrapping C functions) can fail
- In C, we'll often see failure represented as a negative value, with `errno` (see `man errno`) set appropriately (`perror` will print the corresponding error message)
  - Basically every function today can fail in this manner
- In 107, we've mostly ignored this up until this point, but there are a few ways to handle this in C
  - Explicitly check every return value that might fail, write out the failure condition
  - Wrap functions in safe forms (e.g., the textbook creates `Write` from `write`)
  - Use macros to help simplify
  - `gotos` are often used for clean up, but given their potential for misuse, we won't cover them to closely here
  - On the (optional) sockets assignment, we'll provide a few options for error handling (which you should be doing)

# Detour: man pages

- While, in general, we like to tell you to read the manpage for the functions, the man pages for sockets programming tend to be comparatively more difficult to actually find and understand
- I would recommend using the fake man pages from <https://beej.us/guide/bgnet/> and then consulting the real man pages later, as appropriate (and if necessary)

# bind( )

- “bind”s a socket to a particular address/port combo
- `int bind(int sockfd, struct sockaddr *my_addr, int addrlen);`
- Note, we tend to only use `bind` as a server (as a client, we tend not to actually care what our port is)

# struct sockaddr

- `struct sockaddr` is the generic type for a socket address, but we'll use `struct sockaddr_in` or `struct sockaddr_in6` and cast to a `struct sockaddr`

```
struct sockaddr {
    unsigned short sa_family; // address family, AF_xxx
    char sa_data[14]; // 14 bytes of protocol address
};
struct sockaddr_in {
    short int sin_family; // Address family, AF_INET
    unsigned short int sin_port; // Port number
    struct in_addr sin_addr; // Internet address
    unsigned char sin_zero[8]; // Same size as struct sockaddr
};
struct in_addr {
    uint32_t s_addr; // that's a 32-bit int (4 bytes)
};
struct sockaddr_in6 {
    u_int16_t sin6_family; // address family, AF_INET6
    u_int16_t sin6_port; // port number, Network Byte Order
    u_int32_t sin6_flowinfo; // IPv6 flow information
    struct in6_addr sin6_addr; // IPv6 address
    u_int32_t sin6_scope_id; // Scope ID
};
struct in6_addr {
    unsigned char s6_addr[16]; // IPv6 address
};
```

# `inet_pton()`, `inet_addr()`, and `inet_aton()`

- `aton` and `addr` only work for IPv4 addresses
- `int inet_aton(const char *cp, struct in_addr *inp);`
- `in_addr_t inet_addr(const char *cp);`
- `cp` is a string of a dotted quad IP address
- `int inet_pton(int af, const char *src, void *dst);`

# getaddrinfo( )

```
int getaddrinfo(const char *node, // e.g. "www.example.com" or IP
               const char *service, // e.g. "http" or port number
               const struct addrinfo *hints,
               struct addrinfo **res);
```

- Gives us a linked list of struct addrinfos

```
struct addrinfo {
    int ai_flags; // AI_PASSIVE, AI_CANONNAME, etc.
    int ai_family; // AF_INET, AF_INET6, AF_UNSPEC
    int ai_socktype; // SOCK_STREAM, SOCK_DGRAM
    int ai_protocol; // use 0 for "any"
    size_t ai_addrlen; // size of ai_addr in bytes
    struct sockaddr *ai_addr; // struct sockaddr_in or _in6
    char *ai_canonname; // full canonical hostname
    struct addrinfo *ai_next; // linked list, next node
};
```

# bind( )

- `int bind(int sockfd, struct sockaddr *my_addr, int addrlen);`
- Binds our socket to the address and port specified by `my_addr`
- We will often use `INADDR_ANY` to indicate that we want to accept any IPv4 connection (slightly different for IPv6, see “Jumping from IPv4 to IPv6” on Beej’s guide)



# listen( )

- `int listen(int sockfd, int backlog);`
- Starts our socket “listening” (what a server would do)
- `backlog` is how many outstanding requests can be queued until we accept them

# accept ( )

- `int accept(int sockfd, struct sockaddr *addr, socklen_t *addrlen);`
- Returns a file descriptor for a remote connection
- We'll use a `struct sockaddr_storage` (guaranteed large enough to store any address) for the address

```
struct sockaddr_storage {  
    sa_family_t ss_family; // address family  
    // all this is padding, implementation specific, ignore it:  
    char __ss_pad1[_SS_PAD1SIZE];  
    int64_t __ss_align;  
    char __ss_pad2[_SS_PAD2SIZE];  
};
```

# connect ( )

- `int connect(int sockfd, struct sockaddr *serv_addr, int addrlen);`
- Useful for the client, connects our local socket to the remote address

# send ( )

- `int send(int sockfd, const void *msg, int len, int flags);`
- Returns how many bytes were actually sent (may be less than we requested, which we'll have to handle)
- `flags` can be 0 by default
- Note that while we could use `write`, we tend to use `send` instead since it lets us to more specific socket things (see the man page for `flags`)

# recv( )

- `int recv(int sockfd, void *buf, int len, int flags);`
- Returns how many bytes were received (no more than `len`)
- Returns `<0` on error, `0` when remote side has closed

# close ( )

- `close ( sockfd ) ;`
- Prevents any further reads or writes to the socket, the remote peer will receive an error on trying to read or write
- Also, marks the fd as usable again (no longer counts toward our per-process limit)

# shutdown ( )

- `int shutdown(int sockfd, int how);`
- Note that you will still have to `close` eventually

how	Effect
0	Further receives are disallowed
1	Further sends are disallowed
2	Further sends and receives are disallowed (like <code>close()</code> )

**CODE DEMO**



# Handling multiple clients

- We may not get to this in lecture, but you should investigate using `select ( )` and/or `poll ( )` (or `epoll` if you want to get really fancy) for the assignment