# **CS110 Practice Midterm 4 Solution**

This is a closed book, closed note, closed computer exam (although you can use your single double-sided cheat sheet). You have 90 minutes to complete all problems. You don't need to **#include** any header files, and you needn't guard against any errors unless specifically instructed to do so.

Good luck!			
SUNet ID (username):	root@stanford.edu		
Last Name:			
First Name:			
I accept the letter and spirit of the honor code.			
[signed]			
		Score	Grader
1. Game Simulation	on [10]		
2. Parallel <b>grep</b>	[10]		
3. Short Answer C	Questions [10]		
Total	[30]		

#### **Relevant Prototypes**

```
// filesystem access
ssize_t write(int fd, char buffer[], size_t len); // returns num written
int close(int fd); // ignore retval
int pipe(int fds[]); // argument should be array of length 2, ignore retval
int pipe2(int fds[], int flags); // common flag: O_CLOEXEC
int dup2(int old, int new); // ignore retval
#define STDIN FILENO 0
#define STDOUT FILENO 1
#define STDERR FILENO 2
// exceptional control flow and multiprocessing
pid t fork();
pid_t waitpid(pid_t pid, int *status, int flags);
typedef void (*sighandler t)(int sig);
sighandler_t signal(int signum, sighandler_t handler); // ignore retval
int sigsuspend(const sigset_t *mask); // ignore retval
int sigprocmask(int how, const sigset_t *set, sigset_t *oldset); // ign. retval
int execvp(const char *path, char *argv[]); // ignore retval
int kill(pid t pid, int sig); // ignore retval
int setpgid(pid t pid, pid t pgid); // ignore retval
#define WIFEXITED(status) // macro
#define WIFSTOPPED(status) // macro
#define WEXITSTATUS(status) // macro
```

### **Problem 1: Game Simulation [10 points]**

There are artificial intelligence simulators for many games, including tic-tac-toe, checkers, chess, go, etc. Often, an AI improves by playing games against other simulators (or even against itself). For this problem, you will write a program, game\_sim, that runs two game-playing programs and pits them against one another. Your simulator will pass stdout from the first program to stdin of the second program, and vice-versa until both programs terminate. In other words, the first program will produce a move and print it to its stdout, and the second program will expect its std be the first move. Then the second program will print out its move to stdout, and the first program will expect it. This continues until the game terminates. The first game will be responsible for making the first move and will also be responsible for printing the output of the game to the terminal (to stderr, which you can ignore). The way you will run the simulation will be as follows:

```
./game_sim ./game_engine1 ./game_engine2
```

In your program, you will set up the appropriate pipes, and launch the first program (using execvp) as player 1 with the following command line (for this example):

```
./game_engine1
```

The second program should be launched with the following command line:

```
./game_engine2 p2
```

where "p2" tells the program it will be player 2 and will expect the first input from player 1.

The output of the game play will be sent by player 1 to STDERR, and you do not need to worry about this output. Please add your code under the "your code here" section.

```
// file: game-sim.c
#include<stdio.h>
#include<unistd.h> // fork, getpid, getppid
#include <sys/wait.h> // waitpid
#include <fcntl.h> // O CLOEXEC
#include <unistd.h> // pipe, pipe2
int main(int argc, char **argv)
    if (argc < 3) {
        printf("Usage:\n\t %s ./game_engine1 ./game_engine2\n",argv[0]);
        return 0;
    char *game_engine1 = argv[1];
    char *game engine2 = argv[2];
    // your code here:
    int fdsA[2];
    int fdsB[2];
    pipe2(fdsA,O_CLOEXEC);
   pipe2(fdsB,O_CLOEXEC);
   pid_t pids[2];
   pids[0] = fork();
    if (pids[0] == 0) { // child
        dup2(fdsA[1],STDOUT_FILENO);
        dup2(fdsB[0],STDIN_FILENO);
        char *args[3];
        args[0] = game_engine1;
        args[1] = NULL;
        execvp(args[0],args);
    }
   pids[1] = fork();
    if (pids[1] == 0) { // child
        dup2(fdsB[1],STDOUT_FILENO);
        dup2(fdsA[0],STDIN_FILENO);
        char *args[2];
        args[0] = game_engine2;
        args[1] = (char *) "p2";
        args[2] = NULL;
        execvp(args[0],args);
    }
    close(fdsA[0]);
    close(fdsA[1]);
    close(fdsB[0]);
    close(fdsB[1]);
    for (int i = 0; i < 2; i++) {
        waitpid(pids[i],NULL,0);
    }
```

```
return 0;
}
Problem 2: Parallel grep [10 points]
```

The standard UNIX grep program searches lines of STDIN for matching substrings, and prints out matching lines, in order. For example, the following pipeline would search for the substring "cs110" in the printf string, and it would print out three lines:

printf "I love CS110\nCS110 is my favorite\nI like CS107 more\nI have no comment about CS110.\n"  $\mid$  grep CS110

## Output:

```
I love CS110
CS110 is my favorite
I have no comment about CS 110.
```

The interesting part about grep is that it does not, by default, search across line boundaries, meaning that each line search is independent of the others. We can therefore parallelize it using multiprocessing.

Modify the partially written program on the next page that performs the search in parallel. You need to write code at the end of the **printStoppedProcesses** function, and you also need to write the **childHandler** function. Be sure to read through the entire starter code to understand the program flow.

```
const int kMaxProcesses = 80;
int childrenStopped = 0; // global
void printStoppedProcesses(vector<pid t> &children) {
   // wait for all returning jobs, in order, and
   // continue each job to print out its line
   sigset_t suspendset, childset;
   sigemptyset(&suspendset);
   sigemptyset(&childset);
   sigaddset(&childset,SIGCHLD);
   // block SIGCHLD for while loop
   sigprocmask(SIG_BLOCK, &childset, NULL);
   while (children.size() > childrenStopped) {
      sigsuspend(&suspendset);
   sigprocmask(SIG_UNBLOCK, &childset, NULL);
   // your code below
   for (pid_t pid : children) {
        sigprocmask(SIG_BLOCK, &childset, NULL);
        kill(pid,SIGCONT);
        // suspend until exited
        int prevChildrenStopped = childrenStopped;
        while (prevChildrenStopped == childrenStopped) {
           sigsuspend(&suspendset);
        sigprocmask(SIG_UNBLOCK, &childset, NULL);
    }
   // note: blocking could happen outside for loop instead of inside loop.
```

}

```
void childHandler(int sig) {
   // update childrenStopped each time a child stops or exits
   // your code below
    int status;
    while (true) {
        pid_t pid = waitpid(-1, &status, WNOHANG | WUNTRACED);
        if (pid <= 0) break;</pre>
        if (WIFSTOPPED(status)) {
            childrenStopped++;
        if (WIFEXITED(status)) {
            childrenStopped--;
    }
}
void findAndPrint(string line, string searchStr) {
   size_t result = line.find(searchStr);
   raise(SIGSTOP);
   // only print if there is a match
   if (result != string::npos) cout << line << endl;</pre>
int main(int argc, char **argv) {
   if (argc < 2) {
      cout << "usage:\n\t" << argv[0] << " searchstring" << endl;</pre>
      return 1;
   string searchStr = argv[1];
   signal(SIGCHLD, childHandler);
   string line;
   vector<pid_t> children;
   while (getline(cin,line)) {
      if (children.size() >= kMaxProcesses) {
         printStoppedProcesses(children);
         children.clear();
      }
      pid_t pid = fork();
      if (pid == 0) { // child
         findAndPrint(line, searchStr);
         close(STDIN_FILENO); // child should not be reading data
         exit(0);
      children.push_back(pid);
    printStoppedProcesses(children);
    return 0;
}
```

## **Problem 3: Short Answer Questions [10 points]**

Unless otherwise noted, your answers to the following questions should be 75 words or fewer. You needn't write in complete sentences provided it's clear what you're saying. Full credit will only be given to the best of responses. Just because everything you write is true doesn't mean you get all the points.

a. [2 points] The **rename** system call renames a file, moving it from one directory to another if necessary. It comes with the following prototype:

```
int rename(const char *ep, const char *np);
```

**ep** is short for **e**xisting **p**ath, and we'll assume it's an absolute path to a valid file you have permission to rename. **np** (short for **n**ew **p**ath, and also absolute) identifies where the file should be moved to and what new name it should assume. Any intermediate directories needed for the move are created. So, a call to

```
rename("/WWW/index.html","/archive/winter-2019/index-w19.html");
```

would remove **index.html** from **WWW** and move it to **archive/winter-2019**, creating **archive** and **winter-2019** if necessary, with the name of **index-w19.html**. The renaming works even if the file being moved is a directory or a symbolic link.

Without worrying about error checking, describe how **rename** could be efficiently implemented in terms of your Assignment 2 file system.

Find the inode number of the file being moved by iteratively drilling toward it as **pathname\_lookup** would. Remove file's **dirent** from parent directory payload, and then drill toward new parent directory of second argument, creating new **dirent**'s along the way as needed. Finally, append new **dirent** on behalf of new name, with new name (e.g. **index-w19.html**) and existing inode number.

b. [2 points] Consider the following program, which creates and opens a file in read-write mode, writes five characters to the file, and then attempts to read a single character through the same descriptor:

```
int main(int argc, char *argv[]) {
   int f = open("foo.txt", O_RDWR | O_CREAT | O_EXCL, 0644);
   write(f, "abcde", 5);
   int ch;
   int count = read(f, &ch, 1);
   if (count == 0)
      printf("No data!\n");
   else
      printf("Got this: %c\n", ch);
   close(f);
   return 0;
}
```

When the above program runs, the output is **No data!**. Given the output, and leveraging what you know about descriptor, open file, and vnode tables, explain why the output is what it is.

Open file table entry contains a single cursor, which is left to reference the end of the file by the one **write** operation. That means the **read** call attempts to pull from the end of the file, so it returns 0.

c. [2 points] Some students questioned the need to use **pipe2** and **O\_CLOEXEC** while implementing **subprocess**, claiming it was fine to just call **pipe** and manually close all pipe endpoints in the child process before the **execup** call. That, as it turns out, is not true.

Consider the following test program:

```
int main(int argc, char *argv[]) {
   char *args[] = {const_cast<char *>("/usr/bin/sort"), NULL};
   subprocess_t sp1 = subprocess(args, true, false);
   subprocess_t sp2 = subprocess(args, true, false);
   dprintf(sp1.supplyfd, "hello\n");
   dprintf(sp1.supplyfd, "goodbye\n");
   dprintf(sp2.supplyfd, "bonjour\n");
   dprintf(sp2.supplyfd, "aurevoir\n");
   close(sp1.supplyfd);
   waitpid(sp1.pid, NULL, 0);
   close(sp2.supplyfd);
   waitpid(sp2.pid, NULL, 0);
   return 0;
}
```

When the above program is run with a **subprocess** implementation that uses **pipe2** and **O\_CLOEXEC**, we see the expected output. When the above program is run using a **subprocess** implementation that just uses **pipe** (and manually closes the pipe endpoints in the child, before **execvp**), the process hangs.

Where does the program hang? And why?

Program hangs within the first waitpid call. Why? Because spl.supplyfd is active when the second subprocess call (and its call to fork) executes, and that fork clones spl.supplyfd and promotes its reference count by 1. subprocess doesn't know to close it, and spl.supplyfd wasn't created with O\_CLOEXEC. That means the first child running sort never gets an EOF, because spl.supplyfd's reference count never hits 0.

- d. [2 points] When **fork** is called, the set of currently blocked signals is preserved, and the set of installed signal handlers is preserved as well. When **execvp** is called, the set of currently blocked signals is preserved, but all installed signal handlers are cleared and replaced by default handlers.
  - Defend the decision to preserve blocked signal sets across **execvp** boundaries.

Why? The decision to preserve signal masks is consistent with the decision to leave other process-oriented data (e.g. file descriptors, pid, etc) intact, even across **execvp** boundaries.

• Explain why **execvp** restores all signal handlers to be the defaults.

It has no choice, because the code for the original signal handlers is overwritten by a new text segment.

e. [2 points] Consider the following **restartJob** function from my own **stsh** solution, where all error checking has been removed.

```
static void restartJob(const command& command, STSHJobState state) {
    size_t num = parseNumber(command.tokens[0]);
    STSHJob& job = joblist.getJob(num);
    forwardSignal(job, SIGCONT);
    job.setState(state);
    if (state == kForeground)
        waitForForegroundJob(job);
}
```

Because **restartJob** accesses and updates the global job list, one or more signals need to be blocked while **restartJob** is executing. Since you've only installed handlers for **SIGINT**, **SIGTSTP**, and **SIGCHLD**, those are the only three candidates. Which signal or signals need to be blocked, and why?

**SIGCHLD** absolutely needs to be blocked, since it does surgery on the **joblist**, possibly in the middle of the **joblist.getJob(num)** call above. **SIGINT** and **SIGTSTP** should be blocked too, since their handlers would behave differently depending on whether **job.setState(state)** had executed at the time **SIGINT** or **SIGTSTP** was delivered.