Class 9

Compressed Sensing and the RIP and a recap of what we've seen so far

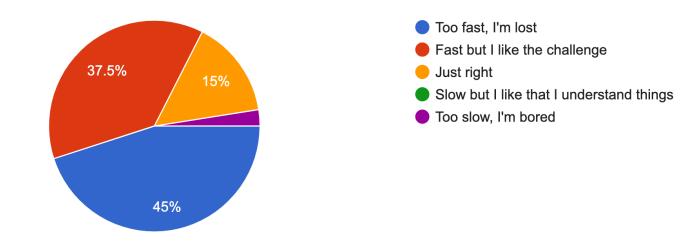
Announcements

- HW4 due Friday!
- I won't be here Thursday, but you will be in capable hands.
 - Prasanna will lead class!

Results from feedback form

Main piece of feedback:

How are you finding the pace of class? 40 responses



- I will try to slow down in class!
 - For today's class, we'll go back to go over some previous material folks asked for.

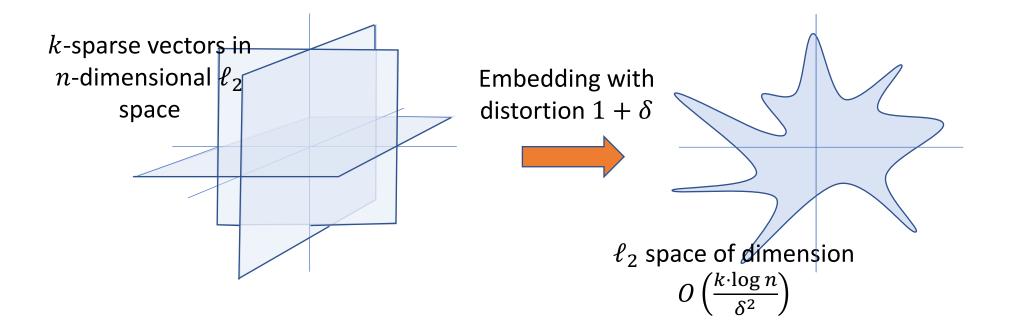
Other things that came up frequently

- The class/HW takes a lot of time.
 - Noted! HW5 should be a bit lighter.
- The quizzes take a lot of time.
 - That is not our intention! If you are struggling with the quizzes, PLEASE ask about them in OH or on Ed (privately). You can also discuss them with your HW group.

- Can we have more practice problems?
 - Check out the problems in "Probability and Computing" by Mitzenmacher and Upfal. The course staff are happy to chat about them with you.
 - We'll have a few more today ©

Note: there were several other things that came up, but not frequently. We did read them and we'll do our best!

Recap



- Restricted Isometry Property: Embedding sparse vectors in lowdimensional space.
- A matrix $A \in \mathbb{R}^{m \times n}$ has the **Restricted Isometry Property** (RIP) with parameters k, δ if for any $x \in \mathbb{R}^n$ with $||x||_0 \le k$,

$$(1 - \delta)||x||_2 \le ||Ax||_2 \le (1 + \delta)||x||_2$$

• Random Gaussian Matrices have the RIP whp! Application: "Compressed Sensing"

Questions?

Minilectures, quiz?

Q1 Does it have the RIP?

6 Points

Which of the following matrices have the RIP with high probability (say, at least 0.9)?

Below, if you check "yes," the matrix should have the (k,ϵ) -RIP with high probability for all sufficiently large n; for all k with $k=\omega(1)$ and $k\leq \sqrt{n}$; and for $\epsilon=0.1$.

Q1.1

2 Points

The identity matrix $I \in \mathbb{R}^{n \times n}$.

- Yes, this has the RIP with high probability.
- O No, this does not have the RIP with high probability.

Q1.2

2 Points

A matrix $A\in\mathbb{R}^{m imes n}$ so that each entry of A is an independent N(0,1/m) random variable, where $m\geq Ck\log n$, for an appropriately large constant C.

- $oldsymbol{\odot}$ Yes, this has the RIP with high probability. (For large enough C).
- igcirc No, this does not have the RIP with high probability. (No matter how large C is).

Q1.3

1 Point

The matrix A from part 1.2, except that one of the columns (say, the first one) has been replaced with all zeros.

- O Yes, this has the RIP with high probability.
- No, this does not have the RIP with high probability.

Q1.4

1 Point

The matrix A from part 1.2, except that one of the rows (say, the first one) has been replaced with all zeros.

- Yes, this has the RIP with high probability.
- O No, this does not have the RIP with high probability.

Q2

6 Points

Let $B^n_p(
ho)=\{x\in\mathbb{R}^n\,:\,\|x\|_p\leq
ho\}$ denote the ℓ_p ball of radius ho in \mathbb{R}^n .

Suppose that $0<\epsilon<1/4$.

Which of the following is a true statement?

Q2.1

3 Points

You can cover $B_2^n(1)$ with $(3/\epsilon)^n$ copies (appropriately shifted) of $B_2^n(\epsilon)$.

- True
- O False

Q2.2

3 Points

You can cover $B_1^n(1)$ with $(3/\epsilon)^n$ copies (appropriately shifted) of $B_1^n(\epsilon)$.

The difference between this one and Q2.1 is that it's ℓ_1 balls instead of ℓ_2 balls. In case it matters, we have $\operatorname{Vol}(B_1^n(\rho))=C_n'\rho^n$ for some constant C_n' , just like for ℓ_2 balls.

- True
- O False

Plan for today

• Slow down and recap!

More comprehensive recap

OMG it's week 5 already -- what have we seen so far? \leftarrow

ALGORITHMS / APPLICATIONS

- · Polynomial Identity Testing
- · Karger's alg for min cuts
- · Primality testing
- Sampling-based Median algorithm
- · Randomized routing on the hypercube
- · Load balancing/powersf 2 choices
- · Finding (appx) sparsest cuts
- · Approximate Nearest Neighbors

TECHNIQUES

- · Linearity of 15
- Writing things as sums of indicator n.v.'s
 Markov's inequality
- · Chebysher's inequality
- · Chemoff bounds + MGFs
- · "Poissonization"

METRIC EMBEDDINGS (both technique and "application")

- Bourgain's embedding
 Locality-Sensitive Hoshing
 Tohnson-Lindenstrauss lemma
 Compressed Sensing / RIP

When to use what bound?

$$\mathbb{P}[X=i] = \frac{i^{-3}}{S(3)} \quad \text{for } i=1,2,...$$

Markov?

Chebysher?

Other?

When to use what bound?

$$\mathbb{P}[X=i] = \frac{i^{-4}}{5(4)}$$

Markov?

Chebysher?

Other?

When to use what bound?

$$\mathbb{P}[X=i] = i^{-k} \leq (k)$$

Markov?

Chebyshov?

Other?

$$S(k) := \sum_{i=1}^{\infty} i^{-k}$$

 $X_i \sim Ber(p)$ iid. $X = \Sigma_i X_i$.

How big does & need to be to show...

When to use what bound?

	Charnof (mult)	Chamoff (add.)	Chebyshow
$P=4$, $P[X=\mu(1+\epsilon)] \leq 0.001$?			
$P=\langle n, P[X=\mu(1+\epsilon)] \leq 0.001$?			
$P = \frac{1}{4} \mathbb{P}[X = \mu(1+\epsilon)] \leq S?$			

 $X_i \sim Ber(p)$ iid. $X = \Sigma_i X_i$.

How big does & need to be to show...

When to use what bound?

	Charnof (mult)	Chamoff (add.)	Chebyshov
$P=4$, $\mathbb{P}[X=\mu(1+\epsilon)] \leq 0.001$?	E = ((Tn)	E= ()(m)	E= (1/m)
$P = \frac{1}{4} \mathbb{P}[X \ge \mu(1+\epsilon)] \le S?$	$\varepsilon = \mathcal{O}\left(\sqrt{\frac{\log(\sqrt{\epsilon})}{n}}\right)$	$\xi = \left(\sqrt{\frac{\log(1/\delta)}{N}} \right)$	$\varepsilon = \left(\frac{1}{\sqrt{\delta n}}\right)$ (much worse?)
$P = \langle n, P[X = \mu(1+\epsilon)] \leq 0.001$?	$\varepsilon = G \left(\frac{1}{n^{1/4}} \right)$	<pre>E = Some constant? Lif it works at all]</pre>	8 = 5 (1/h

Let's practice Poissonization!

How to set up the argument structure

Let's practice Poissonization!

How to set up the argument structure

$$P[Y>0.3n] \leq \frac{\text{Should this}}{\text{be}>n \text{ or } < n?} Y= \Sigma_i Y_i$$
Use Poissonization! Say $k \sim Poi(\frac{1}{2})$, drop k balls into n bins.

n balls into n bins
$$Y_i = 11\{bin i has > 2 balls\}$$