

Review Session 1

Fri. Jan 19, 2:15 pm – 3:05 pm
Thornton 102

1. Agenda

- Announcements
- Homework hints
- MIPS instruction set
- Some PA1 thoughts
- MIPS procedure call convention and example

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2. Announcements

- Homework 1 (can work in groups of 2):
Due date extended to 5 pm on Thursday, January 25.
- The TA email list is ee108b-win0607-tas@lists.stanford.edu
- PA 1 (individual):
Assigned: yesterday
Due 11:59pm on Thursday, February 8, you have 3 weeks.
- Lab 1 (groups of 2):
Due: Tuesday, January 30, you have 12 days.
- Office hours:
Drew: Tuesdays 1pm – 3pm in P129
Yi: Tuesdays 3pm – 5pm P129
- Access to P128 (computer cluster) and P129. Email the TA list for login info.

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3. Homework hints

Problem 1:

- Part b: you can assume all 5000 elements are unequal.
- Part c:
Total Execution Time = (total clk cyc.)/(clk cyc. rate)
- Part d: How does new instruction change the inner loop?
Reduces instruction count, but cycle time has increased...

Problem 2:

- Go line by line of the C code and translate into assembly.
- Tail recursion.

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Problem 3:

- Similar to problem 1.
- No more than 5 bugs.

Problem 4:

- Pseudo instructions.
- Small has 16 bits and Big has 32 bits
- Break big numbers into upper and lower halves
- You may use the notation `upper(big)` and `lower(big)` to extract a portion of a big number.

Please consult the bulletin board first for homework questions, email to the TA list if your question is not already answered.

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4. MIPS instruction set

- R type instruction for arithmetic and logical operations.
The fields are op, rs, rt, rd, shamt, funct
- I type instruction for operations involving immediates up to 16 bits wide. Load, store and branches are formatted as I type.
- J type instructions for jumping.
- Consult the text book's green card or pages 77 and 78 for a comprehensive listing of MIPS instructions.
- Important correction to column 1 of green card:
Shift Left Logical `sll R R[rd]=R[rt]<<shamt`
Shift Right Logical `srl R R[rd]=R[rt]>>shamt`

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5. PA1 Thoughts

- Do the job of a compiler.
- Do not use tail recursion.
- Do use the Partition keyword.
- Know the procedure call conventions well.
- Use display command, xspim or pcspim to debug your code.

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6. MIPS procedure call convention

Call convention summary (please follow this guideline when you do MIPS assembly coding for homeworks, PAs and in exams):

0	\$zero Zero constant (0)	16	\$s0 Callee saved
1	\$at Reserved for assembler	...	
2	\$v0 Expression results	23	\$s7
3	\$v1	24	\$t8 Temporary (cont'd)
4	\$a0 Arguments	25	\$t9
5	\$a1	26	\$k0 Reserved for OS kernel
6	\$a2	27	\$k1
7	\$a3	28	\$gp Pointer to global area
8	\$t0 Caller saved	29	\$sp Stack Pointer
...		30	\$fp Frame Pointer
15	\$t7	31	\$ra Return address

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Example: a recursive program that evaluates factorial 10.

```
main ()
{
    printf ("The factorial of 10 is %d\n", fact (10));
}

int fact (int n)
{
    if (n < 1)
        return (1);
    else
        return (n * fact (n - 1));
}
```

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```
.text                                # Indicate beginning of program
.globl main                          # main can be referenced from other files
main:
    subu    $sp, $sp, 32             # Stack frame is 32 bytes long
    sw     $ra, 20($sp)             # Save return address
    sw     $fp, 16($sp)             # Save old frame pointer
    addiu  $fp, $sp, 28             # Set up frame pointer
    li     $a0, 10                  # Put argument (10) in $a0
    jal   fact                      # Call factorial function
    la    $a0, $LC                  # Load string address in $a0
    move  $a1, $v0                  # Move fact result to $a1
    jal   printf                    # Call the print function
    lw    $ra, 20($sp)             # Restore return address
    lw    $fp, 16($sp)             # Restore frame pointer
    addiu $sp, $sp, 32             # Pop stack frame
    jr    $ra                      # Return to caller

.rdata                                # Read-only data
$LC:
    .ascii "The factorial of 10 is %d\n\000"
    # Store the string in memory, but do not null-terminate it

.text
fact:
    subu    $sp, $sp, 32             # Stack frame is 32 bytes long
    sw     $ra, 20($sp)             # Save return address
    sw     $fp, 16($sp)             # Save frame pointer
    addiu  $fp, $sp, 28             # Set up frame pointer
    sw     $a0, 0($fp)              # Save argument (n)
    lw     $v0, 0($fp)              # Load n
    bgtz  $v0, iterate              # Branch if n > 0
    li    $v0, 1                    # Return 1
    j     pop                       # Jump to code to return

iterate:
    lw     $v1, 0($fp)              # Load n
    subu  $v0, $v1, 1               # Compute n - 1
    move  $a0, $v0                  # Move value to $a0
    jal   fact                      # Call factorial function

    lw     $v1, 0($fp)              # Load n
    mul   $v0, $v0, $v1             # Compute fact(n-1) * n

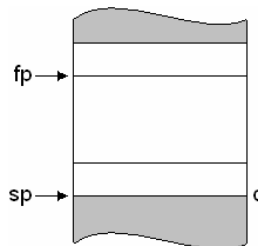
pop:
    lw     $ra, 20($sp)             # Restore $ra
    lw     $fp, 16($sp)             # Restore $fp
    addiu  $sp, $sp, 32             # Pop stack
    jr    $ra                      # Return to caller
```

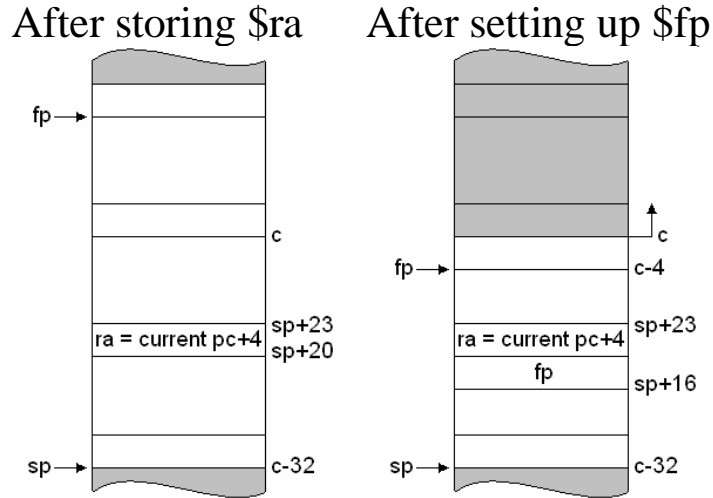
```
iterate:
    lw     $v1, 0($fp)              # Load n
    subu  $v0, $v1, 1               # Compute n - 1
    move  $a0, $v0                  # Move value to $a0
    jal   fact                      # Call factorial function

    lw     $v1, 0($fp)              # Load n
    mul   $v0, $v0, $v1             # Compute fact(n-1) * n

pop:
    lw     $ra, 20($sp)             # Restore $ra
    lw     $fp, 16($sp)             # Restore $fp
    addiu  $sp, $sp, 32             # Pop stack
    jr    $ra                      # Return to caller
```

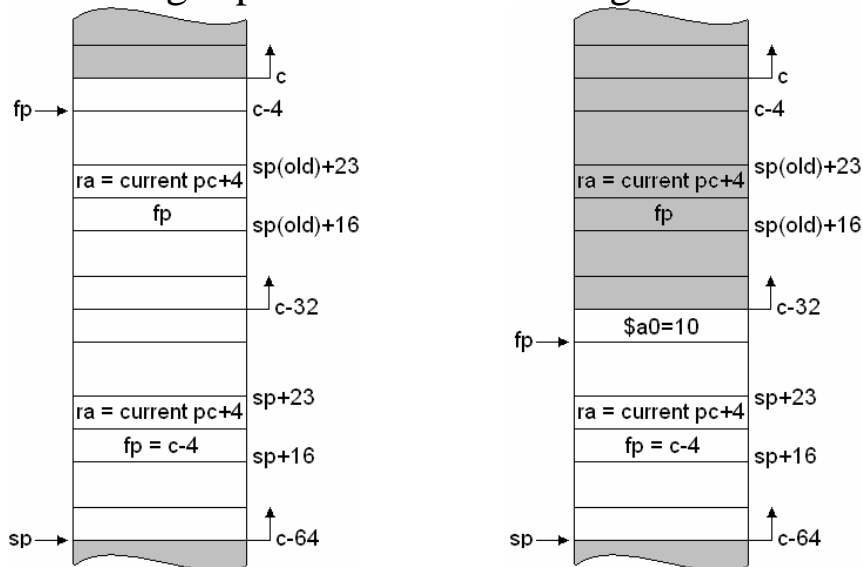
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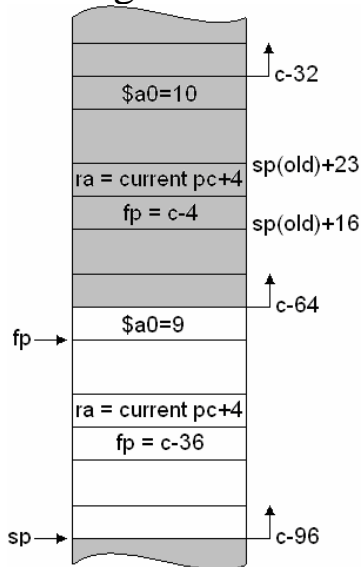
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After saving \$fp in fact After saving \$a0 in iteration 1



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After saving \$a0 in iteration 2



Aside 1:

The `.text` and `.globl` you see in the assembly code are assembler directives, which are there to tell the assembler how to translate a program but which do not translate to machine instructions. Names followed by a colon, such as `str` or `main`, are labels that name the next memory location.

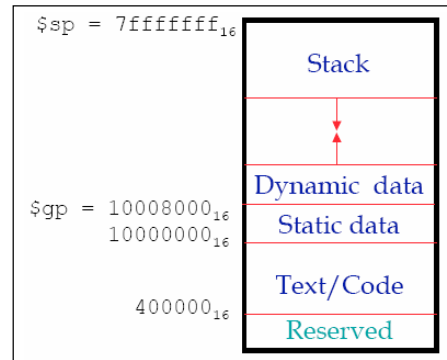
Specifically data definitions start with the `.data` directive, code definition starts with `.text` directive – usually an assembly program would have a `main` as well as a number of subroutine definitions.

Example:

```
.data          # data memory
foo: .word 0   # 32 bit variable

.text         # program memory
.align 2     # word alignment
.globl main  # main is global

main:
    lw $a0, foo
```



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Aside 2:

SPIM provides the “syscall” instruction to communicate with the system and do simple I/O. The specific I/O done depends on the value of some registers.

Example:

```
str:
    .asciiz "jon" # Declare zero-delimited ascii string
    li $v0, 4     # load 4 in $v0 for print_string
                 # syscall
    la $a0, str   # load our string in the syscall
                 # argument register
    syscall      # make the call
```

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In this case the system call code is 4 or `print_string`. Other system calls include:

Code 1: `print_int` with argument `$a0` = integer

Code 2: `print_float` with argument `$f12` = float

Code 3: `print_double` with argument `$f12` = double

Code 4: `print_string` with argument `$a0` = string

Code 5: `read_int` with integer result in `$v0`

Code 6: `read_float` with float result in `$f0`

:
:

The complete listing is available in Appendix A of the textbook (on CD).