Indirect Proofs
Outline for Today

- **What is an Implication?**
  - Understanding a key type of mathematical statement.
- **Negations and their Applications**
  - How do you show something is *not* true?
- **Proof by Contrapositive**
  - What's a contrapositive?
  - And some applications!
- **Proof by Contradiction**
  - The basic method.
  - And some applications!
Logical Implication
Implications

- An *implication* is a statement of the form
  
  If $P$ is true, then $Q$ is true.

- Some examples:
  - Math: If $n$ is an even integer, then $n^2$ is an even integer.
  - Set Theory: If $A \subseteq B$ and $B \subseteq A$, then $A = B$.
  - Queen Bey: If you like it, then you should put a ring on it.
What Implications Mean

• Consider the simple statement

   *If I put fire near cotton, it will burn.*

• Some questions to consider:
  
  • Does this apply to all fire and all cotton, or just some types of fire and some types of cotton? *(Scope)*
  
  • Does the fire cause the cotton to burn, or does the cotton burn for another reason? *(Causality)*

• These are significantly deeper questions than they might seem.

• To mathematically study implications, we need to formalize what implications really mean.
Understanding Implications

“If there's a rainbow in the sky, then it's raining somewhere.”

• In mathematics, implication is directional.
  • The above statement doesn't mean that if it's raining somewhere, there has to be a rainbow.

• In mathematics, implications only say something about the consequent when the antecedent is true.
  • If there's no rainbow, it doesn't mean there's no rain.

• In mathematics, implication says nothing about causality.
  • Rainbows do not cause rain. 😊
What Implications Mean

- In mathematics, a statement of the form
  \textbf{For any } x \textbf{, if } P(x) \textbf{ is true, then } Q(x) \textbf{ is true}
  
  means that any time you find an object \( x \) where \( P(x) \) is true, you will see that \( Q(x) \) is also true (for that same \( x \)).

- There is no discussion of causation here. It simply means that if you find that \( P(x) \) is true, you'll find that \( Q(x) \) is also true.
Implication, Diagrammatically

Set of objects \( x \) where \( P(x) \) is true.

Set of objects \( x \) where \( Q(x) \) is true.

Any time \( P \) is true, \( Q \) is true as well.

If \( P \) isn't true, \( Q \) may or may not be true.
Negations
Negations

• A **proposition** is a statement that is either true or false.
  • Sentences that are questions or commands are not propositions.
• Some examples:
  • If $n$ is an even integer, then $n^2$ is an even integer.
  • $\emptyset = \mathbb{R}$.
  • Moonlight is a good movie.
• The **negation** of a proposition $X$ is a proposition that is true whenever $X$ is false and is false whenever $X$ is true.
• For example, consider the statement “it is snowing outside.”
  • Its negation is “it is not snowing outside.”
  • Its negation is *not* “it is sunny outside.”
How do you find the negation of a statement?
The negation of the *universal* statement

*Every* $P$ *is a* $Q$

is the *existential* statement

*There is a* $P$ *that is not a* $Q$. 
The negation of the \textit{universal} statement

For all \( x \), \( P(x) \) is true.

is the \textit{existential} statement

There exists an \( x \) where \( P(x) \) is false.
The negation of the *existential* statement

There exists a $P$ that is a $Q$

is the *universal* statement

Every $P$ is not a $Q$. 
The negation of the *existential* statement

There exists an $x$ where $P(x)$ is true

is the *universal* statement

For all $x$, $P(x)$ is false.
Puppy Logic

• Consider the statement

I love all puppies.

What is the negation?

A. I don’t love any puppies.
B. I love some puppies and not others.
C. There is at least one puppy I don’t love.

Answer at PollEv.com/cs103 or text CS103 to 22333 once to join, then A, B, or C.
Puppy Logic

- Consider the statement

  **I love all puppies.**

- The following statement is **not** the negation of the original statement:

  **I don’t love any puppies.**

"I love all puppies."

"It's complicated."

"I don't love any puppies."
Puppy Logic

- Consider the statement
  
  **I love all puppies.**

- Here's the proper negation of our initial statement about puppies:
  
  **There's at least one puppy I don't love.**
How do you negate an implication?

Let’s look at:
• Negation of an implication
• A close relative of negation: the Contrapositive
The negation of the statement

“*If* $P$ is true,
*then* $Q$ is true”

is the statement

“$P$ is true,
*and* $Q$ is false.”

*The negation of an implication
is not an implication!*
“If your March Madness bracket is perfect, then you get an A in CS103.”

Which of the following is inconsistent with the above statement?

(A) Your bracket was terrible, and you got an A.
(B) Your bracket was terrible, and you got a B+.
(C) Your bracket was perfect, and you got a B+.
(D) Both (A) and (C)
The negation of the statement

“If your March Madness bracket is perfect, then you get an A in CS103.”

is the statement

“Your March Madness bracket is perfect, and you still didn’t get an A in CS103.

The negation of an implication is not an implication!
The negation of the statement

“For any $x$, if $P(x)$ is true, then $Q(x)$ is true”

is the statement

“There is at least one $x$ where $P(x)$ is true and $Q(x)$ is false.”

*The negation of an implication is not an implication!*
The Contrapositive

- The *contrapositive* of the implication “If $P$, then $Q$” is “If not $Q$, then not $P$.”

- For example:
  - “If your March Madness bracket is perfect, then you get an A in CS103.”
  - Contrapositive: “If you didn’t get an A in CS103 then your March Madness bracket wasn’t perfect.”

- Another example:
  - “If you like it, then you should put a ring on it.”
  - Contrapositive: “If you shouldn’t put a ring on it, then you don’t like it.”
Proof by Contrapositive
To prove the statement

“If \( P \) is true, then \( Q \) is true,”

you could choose to instead prove the equivalent statement

“If \( Q \) is false, then \( P \) is false.”

(if that seems easier).

This is called a \textit{proof by contrapositive}. 
**Theorem:** For any $n \in \mathbb{Z}$, if $n^2$ is even, then $n$ is even.

**Proof:** By contrapositive;

We're starting this proof by telling the reader that it's a proof by contrapositive. This helps cue the reader into what's about to come next.
**Theorem:** For any \( n \in \mathbb{Z} \), if \( n^2 \) is even, then \( n \) is even.

**Proof:** By contrapositive; we prove that if \( n \) is odd, then \( n^2 \) is odd.

Here, we’re explicitly writing out the contrapositive. This tells the reader what we’re going to prove. It also acts as a sanity check by forcing us to write out what we think the contrapositive is.
Theorem: For any $n \in \mathbb{Z}$, if $n^2$ is even, then $n$ is even.

Proof: By contrapositive; we prove that if $n$ is odd, then $n^2$ is odd.

We've said that we're going to prove this new implication, so let's go do it! The rest of this proof will look a lot like a standard direct proof.
**Theorem:** For any \( n \in \mathbb{Z} \), if \( n^2 \) is even, then \( n \) is even.

**Proof:** By contrapositive; we prove that if \( n \) is odd, then \( n^2 \) is odd.

Since \( n \) is odd, there is some integer \( k \) such that \( n = 2k + 1 \). Squaring both sides of this equality and simplifying gives the following:

\[
\begin{align*}
n^2 &= (2k + 1)^2 \\
&= 4k^2 + 4k + 1 \\
&= 2(2k^2 + 2k) + 1.
\end{align*}
\]

From this, we see that there is an integer \( m \) (namely, \( 2k^2 + 2k \)) such that \( n^2 = 2m + 1 \). Therefore, \( n^2 \) is odd. ■
**Theorem:** For any \( n \in \mathbb{Z} \), if \( n^2 \) is even, then \( n \) is even.

**Proof:** By contrapositive; we prove that if \( n \) is odd, then \( n^2 \) is odd.

Since \( n \) is odd, there is some integer \( k \) such that \( n = 2k + 1 \) and similarly ...

The general pattern here is the following:

1. Start by announcing that we’re going to use a proof by contrapositive so that the reader knows what to expect.

2. Explicitly state the contrapositive of what we want to prove.

3. Go prove the contrapositive.
Biconditional

• Combined with what we saw on Wednesday, we have proven that, if \( n \) is an integer:

\[
\text{If } n \text{ is even, then } n^2 \text{ is even.}
\]
\[
\text{If } n^2 \text{ is even, then } n \text{ is even.}
\]

• Therefore, if \( n \) is an integer:

\[
n \text{ is even if and only if } n^2 \text{ is even.}
\]

• “If and only if” is often abbreviated \textit{iff}:

\[
n \text{ is even iff } n^2 \text{ is even.}
\]
Proving Biconditionals

• To prove a theorem of the form $P \iff Q$, you need to prove that $P$ implies $Q$ and that $Q$ implies $P$. (two separate proofs)

• You can use any proof techniques you'd like to show each of these statements.
  - In our case, we used a direct proof for one and a proof by contrapositive for the other.
Proof by Contradiction
“When you have eliminated all which is impossible, then whatever remains, however improbable, must be the truth.”

- Sir Arthur Conan Doyle, *The Adventure of the Blanched Soldier*
Proof by Contradiction

- A *proof by contradiction* is a proof that works as follows:
  - To prove that $P$ is true, assume that $P$ is *not* true.
  - Beginning with this assumption, use logical reasoning to conclude something that is clearly impossible.
    - For example, that $1 = 0$, that $x \in S$ and $x \notin S$, etc.
  - This means that if $P$ is false, something that cannot possibly happen, happens!
  - Therefore, $P$ can't be false, so it must be true.
An Example: *Set Cardinalities*
Set Cardinalities

- We’ve seen sets of many different cardinalities:
  - $|\emptyset| = 0$
  - $|\{1, 2, 3\}| = 3$
  - $|\{n \in \mathbb{N} \mid n < 137\}| = 137$
  - $|\mathbb{N}| = \aleph_0$.
- These span from the finite up through the infinite.

- **Question:** Is there a “largest” set? That is, is there a set that’s bigger than every other set?
**Theorem:** There is no largest set.

**Proof:** Assume for the sake of contradiction that there is a largest set; call it $S$.

To prove this statement by contradiction, we’re going to assume its negation.

What is the negation of the statement “there is no largest set?”

One option: “there is a largest set.”
Theorem: There is no largest set.

Proof: Assume for the sake of contradiction that there is a largest set; call it \( S \).

Notice that we’re announcing

1. that this is a proof by contradiction, and
2. what, specifically, we’re assuming.

This helps the reader understand where we’re going. Remember – proofs are meant to be read by other people!
**Theorem:** There is no largest set.

**Proof:** Assume for the sake of contradiction that there is a largest set; call it $S$.

Now, consider the set $\wp(S)$. By Cantor’s Theorem, we know that $|S| < |\wp(S)|$, so $\wp(S)$ is a larger set than $S$. This contradicts the fact that $S$ is the largest set.

We’ve reached a contradiction, so our assumption must have been wrong. Therefore, there is no largest set. ■
Theorem: There is no largest set.

Proof: Assume for the sake of contradiction that there is a largest set; call it $S$.

The three key pieces:

1. Say that the proof is by contradiction.
2. Say what you are assuming is the negation of the statement to prove.
3. Say you have reached a contradiction and what the contradiction means.

In CS103, please include all these steps in your proofs!

We’ve reached a contradiction, so our assumption must have been wrong. Therefore, there is no largest set. ■
Proving Implications

• To prove the implication

  “If $P$ is true, then $Q$ is true.”

• you can use these three techniques:

  • **Direct Proof.**
    - Assume $P$ and prove $Q$.
  
  • **Proof by Contrapositive.**
    - Assume not $Q$ and prove not $P$.
  
  • **Proof by Contradiction.**
    - ... what does this look like?
Theorem: If \( n \) is an integer and \( n^2 \) is even, then \( n \) is even.

Proof: Assume for the sake of contradiction that ______

What is the assumption?

A. if \( n \) is odd, then \( n^2 \) is odd  
B. \( n \) is an integer and \( n^2 \) is even, and \( n \) is odd  
C. if \( n \) is an integer and \( n^2 \) is odd, then \( n \) is odd  
D. \( n \) is an integer and \( n^2 \) is odd, and \( n \) is odd

Answer at PollEv.com/cs103 or text CS103 to 22333 once to join, then A, B, or C.
**Theorem:** If $n$ is an integer and $n^2$ is even, then $n$ is even.

**Proof:** Assume for the sake of contradiction that $n$ is an integer and that $n^2$ is even, but that $n$ is odd.

Since $n$ is odd we know that there is an integer $k$ such that

$$n = 2k + 1 \quad (1)$$

Squaring both sides of equation (1) and simplifying gives the following:

$$n^2 = (2k + 1)^2 = 4k^2 + 4k + 1 = 2(2k^2 + 2k) + 1 \quad (2)$$

Equation (2) tells us that $n^2$ is odd, which is impossible; by assumption, $n^2$ is even.

We have reached a contradiction, so our assumption must have been incorrect. Thus if $n$ is an integer and $n^2$ is even, then $n$ is even as well. ■
**Theorem:** If \( n \) is an integer and \( n^2 \) is even, then \( n \) is even.  

**Proof:** Assume for the sake of contradiction that \( n \) is an integer and that \( n^2 \) is even, but that \( n \) is odd. Since \( n \) is odd we know that there is an integer \( k \) such that

\[
    n = 2k + 1.  \tag{1}
\]

Squaring both sides of equation (1) and simplifying gives the following:

\[
    n^2 = (2k + 1)^2 \\
    = 4k^2 + 4k + 1 \\
    = 2(2k^2 + 2k) + 1  \tag{2}
\]

Equation (2) tells us that \( n^2 \) is odd, which is impossible; by assumption, \( n^2 \) is even.

We have reached a contradiction, so our assumption must have been incorrect. Thus if \( n \) is an integer and \( n^2 \) is even, \( n \) is even as well. \( \blacksquare \)
**Theorem:** If $n$ is an integer and $n^2$ is even, then $n$ is even.

**Proof:** Assume for the sake of contradiction that $n$ is an integer and that $n^2$ is even, but that $n$ is odd.

Since $n$ is odd we know that there is an integer $k$ such that

$$n = 2k + 1 \quad (1)$$

Squaring both sides of equation (1) and simplifying gives the following:

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Equation (2) tells us that $n^2$ is odd, which is impossible; by assumption, $n^2$ is even.

We have reached a contradiction, so our assumption must have been incorrect. Thus if $n$ is an integer and $n^2$ is even, $n$ is even as well. ■
Recap: Negating Implications

• To prove the statement
  “For any $x$, if $P(x)$ is true, then $Q(x)$ is true”
by contradiction, we do the following:
  • Assume this entire purple statement is false.
  • Derive a contradiction.
  • Conclude that the statement is true.
• What is the negation of the above purple statement?
  “There is an $x$ where $P(x)$ is true and $Q(x)$ is false”
Recap: Contradictions and Implications

• To prove the statement "If $P$ is true, then $Q$ is true" using a proof by contradiction, do the following:
  
  • Assume that $P$ is true and that $Q$ is false.
  • Derive a contradiction.
  • Conclude that if $P$ is true, $Q$ must be as well.
Rational and Irrational Numbers
Rational and Irrational Numbers

- A number $r$ is called a \textit{rational number} if it can be written as

$$r = \frac{p}{q}$$

where $p$ and $q$ are integers and $q \neq 0$.

- A number that is not rational is called \textit{irrational}.
Simplest Forms

• *By definition*, if \( r \) is a rational number, then \( r \) can be written as \( p / q \) where \( p \) and \( q \) are integers and \( q \neq 0 \).

• **Theorem:** If \( r \) is a rational number, then \( r \) can be written as \( p / q \) where \( p \) and \( q \) are integers, \( q \neq 0 \), and \( p \) and \( q \) have no common factors other than 1 and -1.
  • That is, \( r \) can be written as a fraction in simplest form.

• We're just going to take this for granted for now, though with the techniques you'll see later in the quarter you'll be able to prove it!
Question: Are all real numbers rational?
**Theorem:** $\sqrt{2}$ is irrational.

**Proof:** Assume for the sake of contradiction that $\sqrt{2}$ is rational. This means that there must be integers $p$ and $q$ where $q \neq 0$, where $p$ and $q$ have no common divisors other than 1 and -1, and where

$$p / q = \sqrt{2}. \quad (1)$$

Multiplying both sides of equation (1) by $q$ and squaring both sides shows us that

$$p^2 = 2q^2. \quad (2)$$

From equation (2), we see that $p^2$ is even. Earlier, we proved that if $p^2$ is even, then $p$ must also be even. Therefore, we know that there is some integer $k$ such that $p = 2k$. Substituting this into equation (2) and simplifying gives us the following:

$$p^2 = 2q^2$$
$$4k^2 = 2q^2$$
$$2k^2 = q^2 \quad (3)$$

Equation (3) shows that $q^2$ is even. Our earlier theorem tells us that, because $q^2$ is even, $q$ must also be even. But this is not possible - we know that $p$ and $q$ have no common factors other than 1 and -1, but we've shown that $p$ and $q$ must have two as a common factor.

We have reached a contradiction, so our original assumption must have been wrong. Therefore, $\sqrt{2}$ is irrational. $\blacksquare$
**Theorem**: $\sqrt{2}$ is irrational.

**Proof**: Assume for the sake of contradiction that $\sqrt{2}$ is rational. This means that there must be integers $p$ and $q$ where $q \neq 0$, where $p$ and $q$ have no common divisors other than 1 and -1, and where

$$\frac{p}{q} = \sqrt{2}. \quad (1)$$

Multiplying both sides of equation (1) by $q$ and squaring both sides shows us that

$$p^2 = 2q^2. \quad (2)$$

From equation (2), we see that $p^2$ is even. Earlier, we proved that if $p^2$ is even, then $p$ must also be even. Therefore, we know that there is some integer $k$ such that

$$p = 2k. \quad (3)$$

Equation (3) shows that $q^2$ is even. Our earlier theorem tells us that, because $q^2$ is even, $q$ must also be even. But this is not possible - we know that $p$ and $q$ have no common factors other than 1 and -1, but we've shown that $p$ and $q$ must have two as a common factor.

We have reached a contradiction, so our original assumption must have been wrong. Therefore, $\sqrt{2}$ is irrational. □

The three key pieces:

1. State that the proof is by contradiction.
2. State what you are assuming is the negation of the statement to prove.
3. State you have reached a contradiction and what the contradiction entails.

In CS103, please include all these steps in your proofs!
Vi Hart on Pythagoras and the Square Root of Two:

http://www.youtube.com/watch?v=X1E7I7_r3Cw
What We Learned

- **What's an implication?**
  - It's a statement of the form “if $P$, then $Q$,” and states that if $P$ is true, then $Q$ is true.

- **How do you negate formulas?**
  - It depends on the formula. There are nice rules for how to negate universal and existential statements and implications.

- **What is a proof by contrapositive?**
  - It's a proof of an implication that instead proves its contrapositive.
  - (The contrapositive of “if $P$, then $Q$” is “if not $Q$, then not $P$.”)

- **What's a proof by contradiction?**
  - It's a proof of a statement $P$ that works by showing that $P$ cannot be false.
Next Time

- **Mathematical Logic**
  - How do we formalize the reasoning from our proofs?

- **Propositional Logic**
  - Reasoning about simple statements.

- **Propositional Equivalences**
  - Simplifying complex statements.
Handouts

- There are six (!) total handouts for today:
  - Handout 08: Guide to Proofs
  - Handout 09: Mathematical Vocabulary
  - Handout 10: Guide to Indirect Proofs
  - Handout 11: Ten Techniques to Get Unstuck
  - Handout 12: Proofwriting Checklist
  - Handout 13: Problem Set One

- Be sure to read handouts; there's a lot of really important information in there!
Announcements

• Problem Set 1 goes out today!

• *Checkpoint* due Monday, January 15 at 2:30PM.
  • Grade determined by attempt rather than accuracy. It's okay to make mistakes – we want you to give it your best effort, even if you're not completely sure what you have is correct.
  • We will get feedback back to you with comments on your proof technique and style.
  • The more effort you put in, the more you'll get out.

• *Remaining problems* due Friday, January 19 at 2:30PM.
  • Feel free to email staff list with questions, stop by office hours, or ask questions on Piazza!
Submitting Assignments

• All assignments should be submitted through GradeScope.
  • The programming portion of the assignment gets submitted separately from the written component.
  • The written component **must** be typed up; handwritten solutions don’t scan well and get mangled in GradeScope.

• Summary of the late policy:
  • Everyone has *three* 24-hour late days.
  • Late days can't be used on checkpoints.
  • Nothing may be submitted more than *two* days past the due date.
  • Because submission times are recorded automatically, we're strict about the submission deadlines.

• **Very good idea:** Leave at least two hours buffer time for your first assignment submission, just in case something goes wrong.

• **Very bad idea:** Wait until the last minute to submit.
Working in Pairs

• You can work on the problem sets individually or in pairs.

• Each person/pair should only submit a single problem set, and should submit it only once.

• Full details about the problem sets, collaboration policy, and Honor Code can be found in Handout 04 and Handout 05.
A Note on the Honor Code
Office hours have started!

Schedule is available on the course website.
Appendix: Negating Statements
Scoping Implications

- Consider the following statements:
  
  If $A \subseteq B$ and $B \subseteq C$, then $A \subseteq C$.
  
  If $n$ is even, then $n^2$ is even.
  
  If $A \subseteq B$ and $B \subseteq A$, then $A = B$.

- In the above statements, what are $A$, $B$, $C$, and $n$? Are they specific objects? Or do these claims hold for all objects?
Implications and Universals

• In discrete math, most implications involving unknown quantities are, implicitly, universal statements.*

• For example, the statement

   \textbf{If } A \subseteq B \textbf{ and } B \subseteq C, \textbf{ then } A \subseteq C

actually means

   \textbf{For any sets } A, B, \textbf{ and } C, \textbf{ if } A \subseteq B \textbf{ and } B \subseteq C, \textbf{ then } A \subseteq C.

* Your proofs should never use variables without officially introducing them though. This will become more clear next Wednesday.
Negating Universal Statements

“For all $x$, $P(x)$ is true”
becomes
“There is an $x$ where $P(x)$ is false.”

Negating Existential Statements

“There exists an $x$ where $P(x)$ is true”
becomes
“For all $x$, $P(x)$ is false.”

Negating Implications

“For every $x$, if $P(x)$ is true, then $Q(x)$ is true”
becomes
“There is an $x$ where $P(x)$ is true and $Q(x)$ is false”
"If $P(x)$ is true, then $Q(x)$ is true."

"Sometimes $P(x)$ is true and $Q(x)$ is true, -and- sometimes $P(x)$ is true and $Q(x)$ is false."

"$P(x)$ implies $Q(x)$"
\[ q^2 = b + p^2 \]

(Imagine \( q \) is as small as possible.)