

Mathematical Logic

Where We Are Now

- Our coverage of logic focused on these topics:
 - Propositional variables.
 - Propositional connectives.
 - Propositional equivalences.
 - Predicates, functions, and constant symbols.
 - Objects and propositions.
 - Quantifiers.
 - Evaluating first-order formulas relative to a world.
 - Translating into first-order logic.
 - Negating and simplifying first-order formulas.
- Your goal this week is to keep your proof skills sharp while mastering the ins and outs of first-order logic.

Things You Should Do Today

- Review the solutions for Problem Set One and make sure you *completely* and *unambiguously* understand the answers. Ask for help if this isn't the case!
- Read the “Guide to Negating Formulas” and “Guide to First-Order Logic Translations” on the course website to get more exposure and practice with those skills.
- Continue working through PS2.

Things You Should Do Tomorrow

- Look over your feedback on PS1 (and the PS2 checkpoint) and make sure you understand all the feedback you get **completely** and **unambiguously**. Ask the course staff for help, either on Piazza or in office hours, if you don't.
- Continue working on PS2.
- Start reviewing your partner's answers and compiling a single, definitive set of answers that you're going to turn in.
- Stop by office hours to get feedback on your proofs and take that feedback seriously.
- (Also, complete the **CS103A assignment** after Wednesday's lecture—it is due Friday at 2:30 pm!)

Topics You Wanted More Practice With

- First-order logic and translations [today's topic]
- Truth tables
- Proofs and disproofs [more next week]
 - For now, make sure you look at the PS1 solutions and feedback!!

Mechanics: *Negating Statements*

$$\forall p. (Person(p) \rightarrow$$
$$\quad \exists q. (Person(q) \wedge p \neq q \wedge$$
$$\quad \quad Loves(p, q)$$
$$\quad)$$
$$)$$

$$\neg \forall p. (Person(p) \rightarrow$$
$$\exists q. (Person(q) \wedge p \neq q \wedge$$
$$Loves(p, q)$$
$$)$$
$$)$$

$\neg \forall p. (Person(p) \rightarrow$
 $\exists q. (Person(q) \wedge p \neq q \wedge$
 $Loves(p, q)$
)
)

Useful Resource:

Go to cs103.stanford.edu and
read the Guide to Negating
Formulas.

$\neg \forall p. (Person(p) \rightarrow$
 $\exists q. (Person(q) \wedge p \neq q \wedge$
 $Loves(p, q)$
)
)

$$\frac{\neg \forall x. A}{\exists x. \neg A}$$

$\neg \forall p. (Person(p) \rightarrow$
 $\exists q. (Person(q) \wedge p \neq q \wedge$
 $Loves(p, q)$
)
)

$$\frac{\neg \forall x. A}{\exists x. \neg A}$$

$\neg \forall p. (Person(p) \rightarrow$
 $\exists q. (Person(q) \wedge p \neq q \wedge$
 $Loves(p, q)$
)
)

$$\frac{\neg \forall x. A}{\exists x. \neg A}$$

$\exists p. \neg(Person(p) \rightarrow$
 $\exists q. (Person(q) \wedge p \neq q \wedge$
 $Loves(p, q)$
)
)

$$\frac{\neg \forall x. A}{\exists x. \neg A}$$

$$\exists p. \neg(Person(p) \rightarrow$$
$$\exists q. (Person(q) \wedge p \neq q \wedge$$
$$Loves(p, q)$$
$$)$$
$$)$$

$\exists p. \neg(Person(p) \rightarrow$
 $\quad \exists q. (Person(q) \wedge p \neq q \wedge$
 $\quad \quad Loves(p, q)$
 $\quad)$
 $)$

$$\frac{\neg(A \rightarrow B)}{A \wedge \neg B}$$

$\exists p. \neg(Person(p) \rightarrow$
 $\exists q. (Person(q) \wedge p \neq q \wedge$
 $Loves(p, q)$
)
)

$$\frac{\neg(A \rightarrow B)}{A \wedge \neg B}$$

$\exists p. \neg(Person(p) \rightarrow$
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 $)$
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$$)$$

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 $Loves(p, q)$
)
)

$$\frac{\neg \exists x. A}{\forall x. \neg A}$$

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 $\neg \exists q. (Person(q) \wedge p \neq q \wedge$
 $Loves(p, q)$
)
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 $Loves(p, q)$
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$$\frac{\neg \exists x. A}{\forall x. \neg A}$$

$$\exists p. (Person(p) \wedge$$
$$\quad \forall q. \neg(Person(q) \wedge p \neq q \wedge$$
$$\quad \quad Loves(p, q)$$
$$\quad)$$
$$)$$

$\exists p. (Person(p) \wedge$
 $\forall q. \neg(Person(q) \wedge p \neq q \wedge$
 $Loves(p, q)$
)
)

$$\frac{\neg(A \wedge B)}{A \rightarrow \neg B}$$

$\exists p. (Person(p) \wedge$
 $\quad \forall q. \neg(Person(q) \wedge p \neq q \wedge$
 $\quad \quad Loves(p, q)$
 $\quad)$
 $)$

$$\frac{\neg(A \wedge B)}{A \rightarrow \neg B}$$

$\exists p. (Person(p) \wedge$
 $\forall q. \neg(Person(q) \wedge p \neq q \wedge$
 $Loves(p, q)$
)
)

$$\frac{\neg(A \wedge B)}{A \rightarrow \neg B}$$

$\exists p. (Person(p) \wedge$
 $\forall q. (Person(q) \wedge p \neq q \rightarrow$
 $\neg Loves(p, q)$
)
)

$$\frac{\neg(A \wedge B)}{A \rightarrow \neg B}$$

$$\exists p. (Person(p) \wedge$$
$$\quad \forall q. (Person(q) \wedge p \neq q \rightarrow$$
$$\quad \quad \neg Loves(p, q)$$
$$\quad)$$
$$)$$

$$\forall p. (Person(p) \rightarrow$$
$$\exists q. (Person(q) \wedge p \neq q \wedge$$
$$Loves(p, q)$$
$$)$$
$$)$$
$$\exists p. (Person(p) \wedge$$
$$\forall q. (Person(q) \wedge p \neq q \rightarrow$$
$$\neg Loves(p, q)$$
$$)$$
$$)$$

$\exists p. (Person(p) \wedge$
 $\forall q. (Person(q) \wedge p \neq q \rightarrow$
 $Loves(q, p)$
)
)

Your turn!

Try negating this formula with
the other folks at your table.

See what you come up with!

$$\neg \exists p. (Person(p) \wedge$$
$$\forall q. (Person(q) \wedge p \neq q \rightarrow$$
$$Loves(q, p)$$
$$)$$
$$)$$

$$\forall p. \neg(Person(p) \wedge$$
$$\forall q. (Person(q) \wedge p \neq q \rightarrow$$
$$Loves(q, p)$$
$$)$$
$$)$$

$$\forall p. (Person(p) \rightarrow$$
$$\neg \forall q. (Person(q) \wedge p \neq q \rightarrow$$
$$Loves(q, p)$$
$$)$$
$$)$$

$$\forall p. (Person(p) \rightarrow$$
$$\quad \exists q. \neg(Person(q) \wedge p \neq q \rightarrow$$
$$\quad \quad Loves(q, p)$$
$$\quad)$$
$$)$$

$$\forall p. (Person(p) \rightarrow$$
$$\quad \exists q. (Person(q) \wedge p \neq q \wedge$$
$$\quad \quad \neg Loves(q, p)$$
$$\quad)$$
$$)$$

$$\begin{aligned} & \exists p. (Person(p) \wedge \\ & \quad \forall q. (Person(q) \wedge p \neq q \rightarrow \\ & \quad \quad Loves(q, p) \\ & \quad) \\ &) \end{aligned}$$
$$\begin{aligned} & \forall p. (Person(p) \rightarrow \\ & \quad \exists q. (Person(q) \wedge p \neq q \wedge \\ & \quad \quad \neg Loves(q, p) \\ & \quad) \\ &) \end{aligned}$$

Techniques: ***Translating Statements***

Common Patterns

- A statement of the form

$$\forall x. (P(x) \rightarrow Q(x))$$

can be read as “all P 's are Q 's.”

- A statement of the form

$$\exists x. (P(x) \wedge Q(x))$$

can be read as “there is a P that is also a Q ” or “some P 's are Q 's.”

- **Remember:** If you see \exists paired with \rightarrow or \forall paired with \wedge , the statement is probably incorrect!

Given the predicates

- $Person(p)$, which states that p is a person, and
- $CanLearnFrom(x, y)$, which says that x can learn from y ,

write a statement in first-order logic that says “everyone has someone they can learn from.”

Everyone has someone they can learn from

Every person p has someone they can learn from

Every person p has someone they can learn from

“All A s are B s.”

$\forall x. (A(x) \rightarrow B(x))$

$\forall p. (Person(p) \rightarrow$
 p has someone they can learn from
)

$\forall p. (Person(p) \rightarrow$
there is a person q that p can learn from
)

$\forall p. (\text{Person}(p) \rightarrow$

there is a person q that p can learn from

)

“Some As are Bs.”

$\exists x. (A(x) \wedge B(x))$

$\forall p. (Person(p) \rightarrow$
 $\exists q. (Person(q) \wedge$
 p can learn from q
)
)

$$\forall p. (Person(p) \rightarrow$$
$$\quad \exists q. (Person(q) \wedge$$
$$\quad \quad CanLearnFrom(p, q)$$
$$\quad)$$
$$)$$

$$\forall p. (Person(p) \rightarrow$$
$$\quad \exists q. (Person(q) \wedge$$
$$\quad \quad CanLearnFrom(p, q)$$
$$\quad)$$
$$)$$

- Pair quantifiers with the appropriate connectives.*
- Use whitespace and indentation to clarify meaning.*
- Defer quantifiers until they're needed.*
- Check the types of all terms in the formula.*
- Check the scoping on each variable.*

Consider this statement:

“If someone is happy, then everyone is happy.”

What is the *contrapositive* of this statement?

If someone is happy, then everyone is happy

someone is happy → *everyone is happy*

someone is happy $\rightarrow (\forall x. \text{Happy}(x))$

$(\exists x. \textit{Happy}(x)) \rightarrow (\forall x. \textit{Happy}(x))$

$(\exists x. \textit{Happy}(x)) \rightarrow (\forall x. \textit{Happy}(x))$

$(\exists x. \textit{Happy}(x)) \rightarrow (\forall x. \textit{Happy}(x))$

$$\neg(\forall x. \text{Happy}(x)) \rightarrow \neg(\exists x. \text{Happy}(x))$$

$(\exists x. \neg \text{Happy}(x)) \rightarrow \neg(\exists x. \text{Happy}(x))$

$$(\exists x. \neg \text{Happy}(x)) \rightarrow (\forall x. \neg \text{Happy}(x))$$

CONTRAPOSITIVE of “If someone is happy, then everyone is happy”

$$(\exists x. \neg \text{Happy}(x)) \rightarrow (\forall x. \neg \text{Happy}(x))$$

“If someone is not happy, then everyone is not happy.”

Consider this statement:

“If someone is happy, then everyone is happy.”

What is the *negation* of this statement?

$(\exists x. \textit{Happy}(x)) \rightarrow (\forall x. \textit{Happy}(x))$

$(\exists x. \text{Happy}(x)) \rightarrow (\forall x. \text{Happy}(x))$

$\neg((\exists x. \text{Happy}(x)) \rightarrow (\forall x. \text{Happy}(x)))$

$(\exists x. \text{Happy}(x)) \rightarrow (\forall x. \text{Happy}(x))$

$(\exists x. \text{Happy}(x)) \wedge \neg(\forall x. \text{Happy}(x))$

$(\exists x. \textit{Happy}(x)) \rightarrow (\forall x. \textit{Happy}(x))$

$(\exists x. \textit{Happy}(x)) \wedge (\exists x. \neg \textit{Happy}(x))$

“If someone is happy, then everyone is happy”

$(\exists x. \text{Happy}(x)) \rightarrow (\forall x. \text{Happy}(x))$

$(\exists x. \text{Happy}(x)) \wedge (\exists x. \neg \text{Happy}(x))$

“Someone is happy and someone is not happy.”