



YEAH A3: Recursion!!!

CS106B Summer '21: Assignment 3
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Assignment 3 ~Logistics~

A3 is due on **Thursday, July 15** at 11:59pm PT

Grace period until **Saturday, July 17** at 11:59pm PT

Note: this assignment has a lot of parts, so start early!!

(Seriously)

recursion



 All

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Tools

About 143,000,000 results (0.40 seconds)

Did you mean: ***recursion***



Before anything, some wise words from **char** the charmander:

Recursion is all about breaking down a problem into smaller and smaller pieces until you reach the simplest case that you can't break down any further.

We should have a strong idea of how to break down the task at hand before even touching the keyboard.

To help, **DRAW LOTS OF PICTURES AS YOU TACKLE THIS ASSIGNMENT!**



Assignment 3

1. Sierpinski
2. Balanced
3. Mergesort
4. Boggle

But first! Warmups!

- 1) Examine Recursive Stack Frames
- 2) Recognize Stack Overflow (kind of like the recursive version of entering an infinite loop)
- 3) Test a Recursive Function
- 4) Debugging a Recursive Function

But first! Warmups!



Assignment 3

1. **Sierpinski**
2. Balanced
3. Mergesort
4. Boggle

Sierpinski

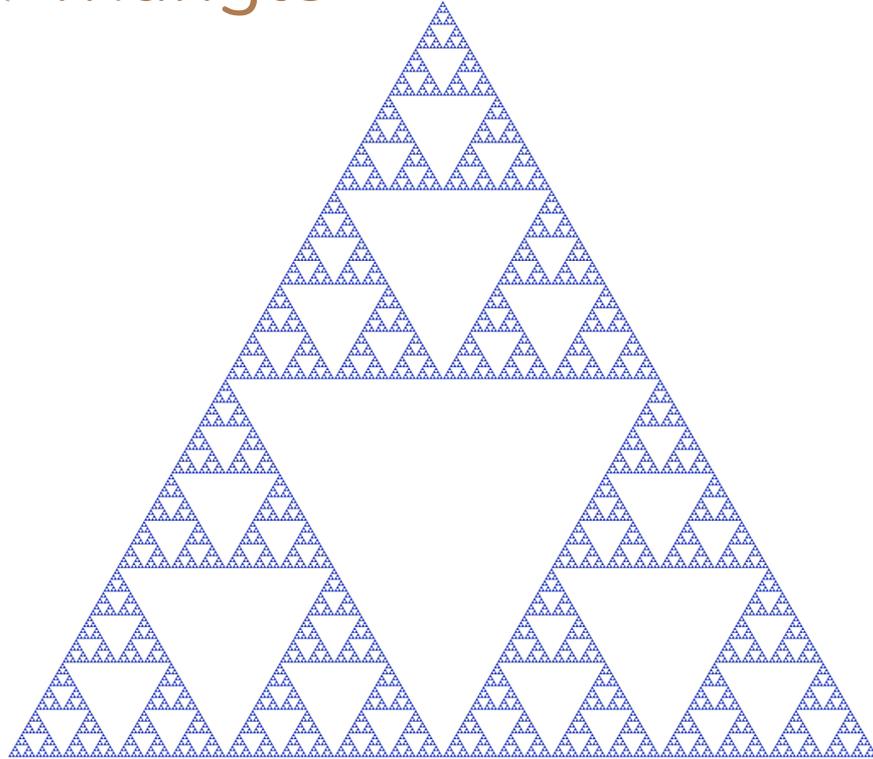
A Polish Mathematician (1882-1969)

Famous for triangles!

Huh?



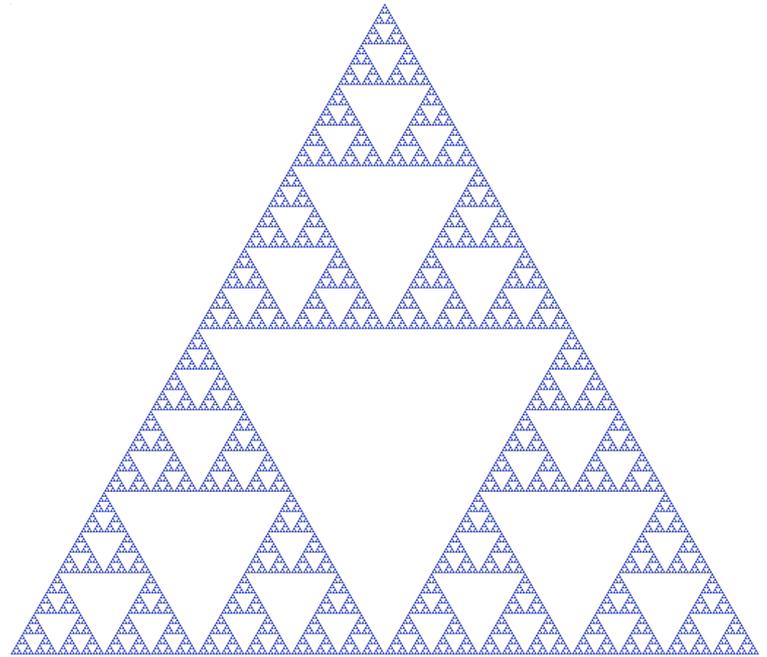
Sierpinski Triangle



Sierpinski Triangle

We can use *recursion* to create incredible graphic designs of **fractals** (images that have self-similar subparts)

Sierpinski Triangle is a famous example of fractals!



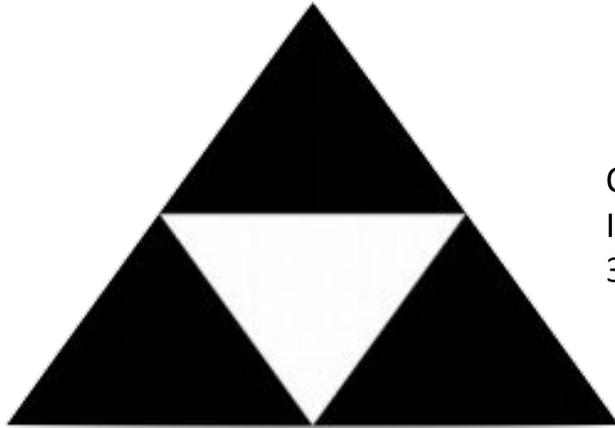
Sierpinski Triangle

Sierpinski Triangles are defined recursively:

- An order-0 Sierpinski triangle is a plain filled triangle.
- An order- n Sierpinski triangle, where $n > 0$, consists of three Sierpinski triangles of order $n - 1$, whose side lengths are half the size of the original side lengths, arranged so that they meet corner-to-corner.

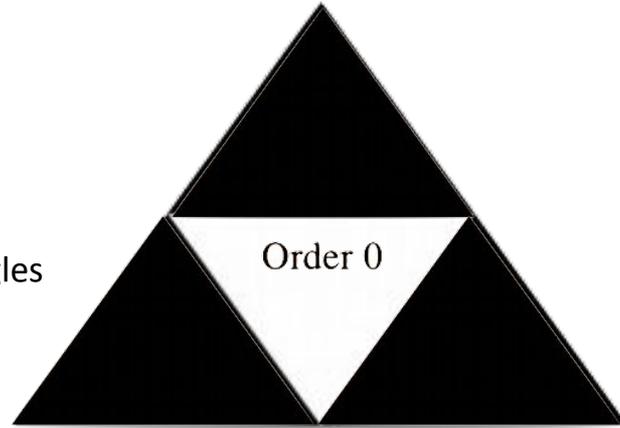


Sierpinski Triangle



Order 1

Order-1
Is the same as:
3 order-0 triangles



Order 0 Order 0
Order 1

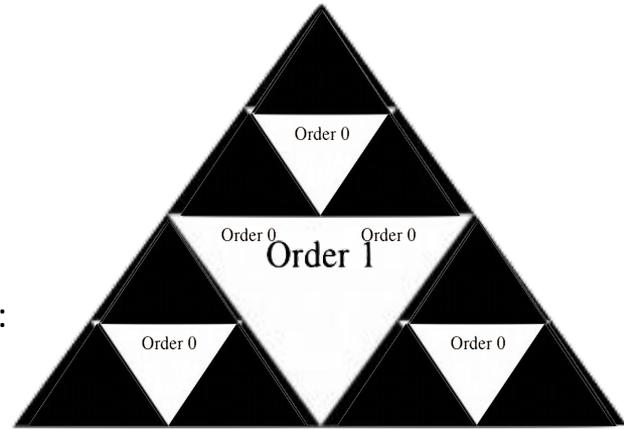
Sierpinski Triangle



Order 2

Order-2
Is the same as:
3 order-1 triangles
Which is the same as:
9 order-0 triangles

And so on...



Order 0 Order 0 Order 0
Order 1
Order 0 Order 0 Order 0
Order 1
Order 2

Sierpinski Triangle

You could say:

- **Drawing a sierpinski triangle of Order N is just like saying:**
 - Drawing a sierpinski triangle of Order N-1
 - Drawing a sierpinski triangle of Order N-1
 - Drawing a sierpinski triangle of Order N-1
- (With different side lengths and positions of course)

Sierpinski Triangle

- **Drawing a sierpinski triangle of Order N is just like saying:**
 - **Drawing a sierpinski triangle of Order N-1**
 - Drawing a sierpinski triangle of Order N-1
 - Drawing a sierpinski triangle of Order N-1
 - Drawing a sierpinski triangle of Order N-1
 - **Drawing a sierpinski triangle of Order N-1**
 - Drawing a sierpinski triangle of Order N-1
 - Drawing a sierpinski triangle of Order N-1
 - Drawing a sierpinski triangle of Order N-1
 - **Drawing a sierpinski triangle of Order N-1**
 - Drawing a sierpinski triangle of Order N-1
 - Drawing a sierpinski triangle of Order N-1
 - Drawing a sierpinski triangle of Order N-1

Sierpinski Triangle

You will be implementing a function:

```
int drawSierpinskiTriangle(GWindow& window, GPoint one, GPoint two, GPoint three, int order)
```



Graphics window that will show the visuals



Three points of a triangle
GPoint is {x, y} coordinates



order number

You will be calculating the number of triangles you draw (recursively!)

Sierpinski Triangle

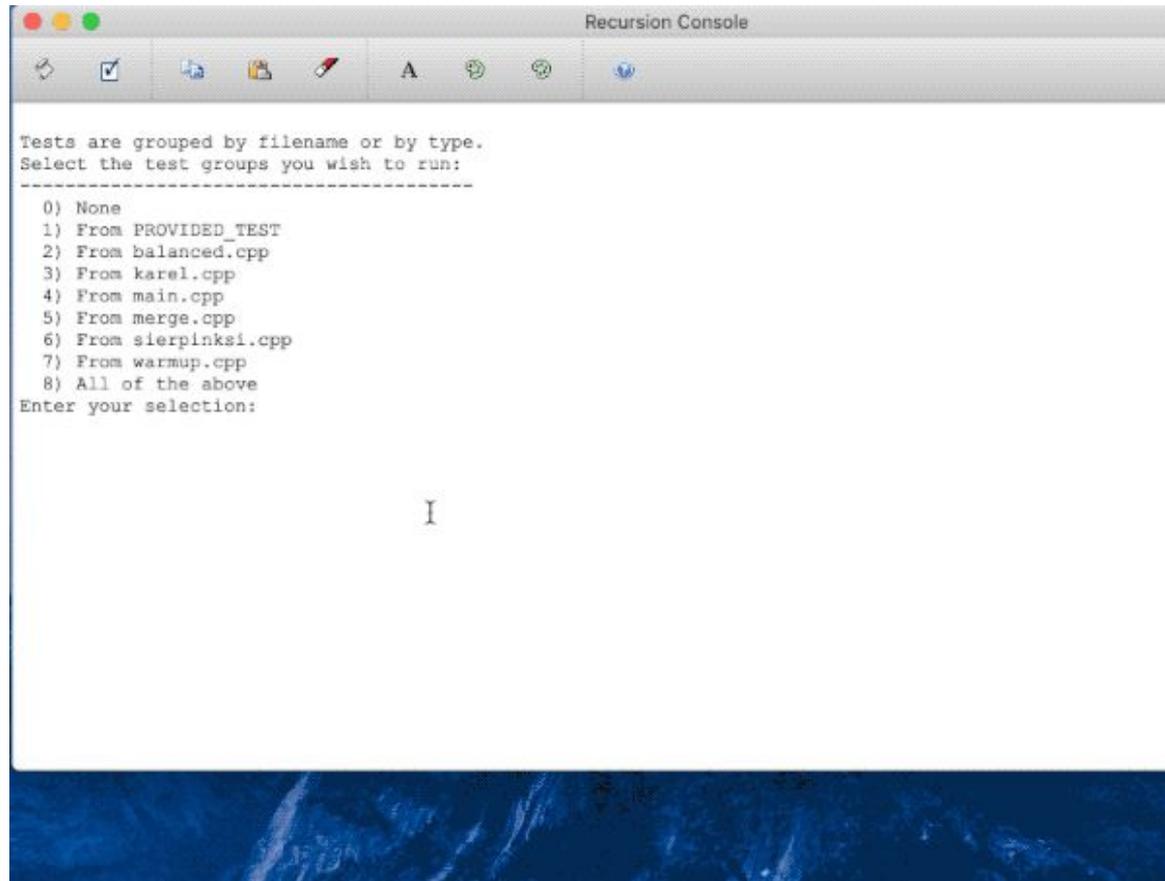
We provide you with the function:

```
void fillBlackTriangle(GWindow& gw, GPoint one, GPoint two, GPoint three)
```

Which draws a triangle onto the graphics window, given the window, and all three points.

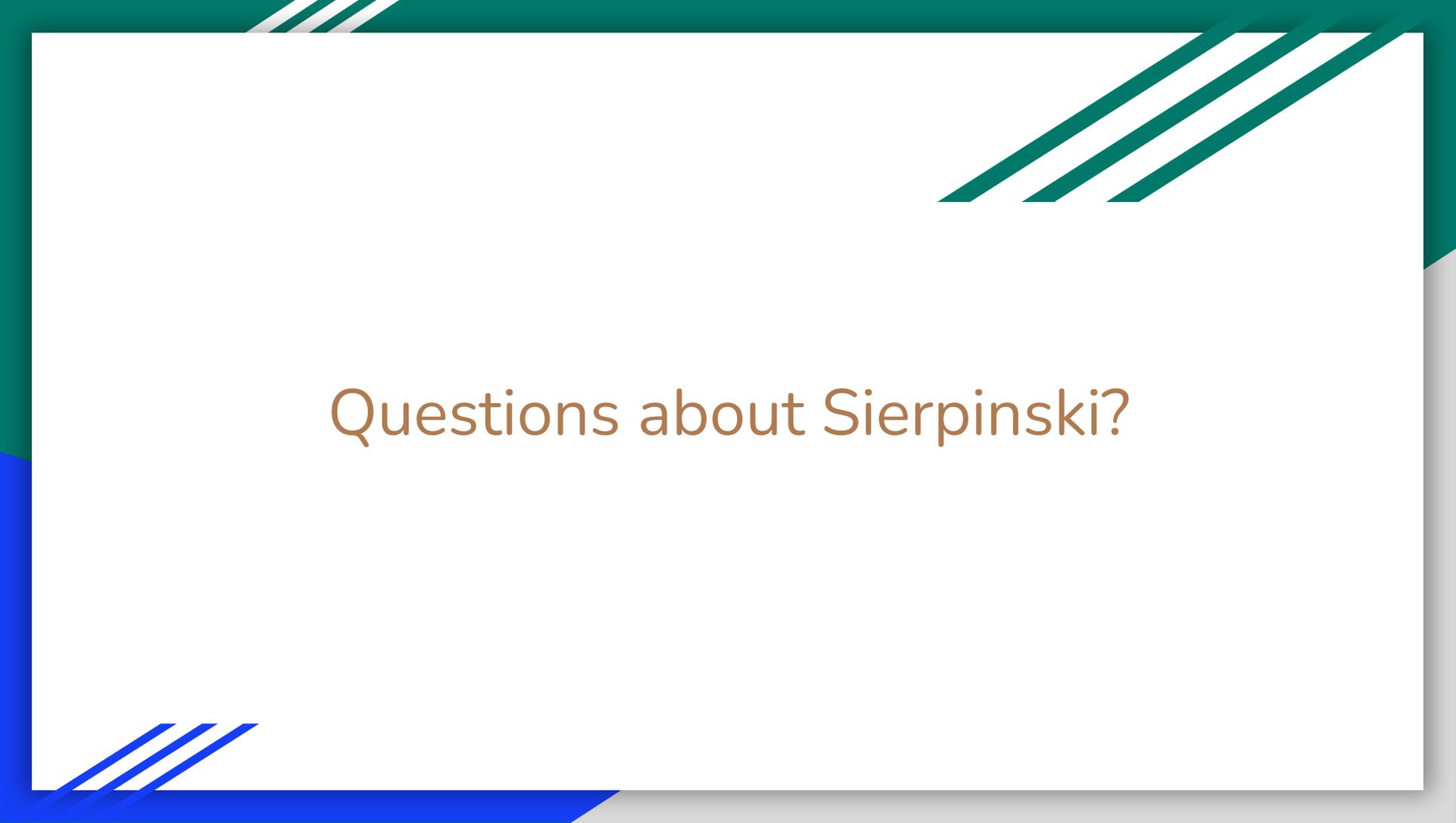
(Hint: you only want to call this function **once** in drawSierpinskiTriangle (base case??))

Trust your recursion to draw each triangle in each correct position!



Sierpinski Triangle Tips/Tricks

- Draw pictures!! This will be SUPER helpful in figuring out the location for each triangle corner!
- The Midpoint of each side of the outer larger triangle becomes a *corner* of one of the smaller inner triangles.
 - Given a line segment whose endpoints are $\{x_1, y_1\}$ and $\{x_2, y_2\}$, its midpoint is located at $\{ (x_1+x_2)/2, (y_1+y_2)/2 \}$
- Your triangles do not need to be equilateral!
- A **GPoint** struct is used to specify a x,y coordinate within the graphics window. Here is the documentation for [GPoint](#).
- The only drawing function you need is our provided **fillBlackTriangle**. You do not need to dig into any other drawing functionality in the Stanford libraries.



Questions about Sierpinski?

Assignment 3

1. Sierpinski
2. **Balanced**
3. Mergesort
4. Boggle



PERFECTLY BALANCED

AS ALL THINGS SHOULD BE

memegenerator.net

Balanced

In most programming languages, we see *bracketing operators* such as:

(. . .) -- Parenthesis (?)

[. . .] -- Square Bracket (?)

{ . . . } -- Curly Bracket (?)

Balanced

Let's look at how they are used!

```
int main() {  
    int x = 2 * (vec[2] + 3); x = (1 + random());  
}
```

Balanced

Let's look at how they are used!

```
int main() {  
    int x = 2 * (vec[2] + 3); x = (1 + random());  
}
```

Operators only → `() { ([]) () }`

Yes, the operators are correctly balanced!

Balanced

What is an “unbalanced” pair?

```
( ( [ a ] )   The line is missing a close parenthesis.  
3 ) (        The close parenthesis comes before the open parenthesis.  
{ ( x } y )  The parentheses and braces are improperly nested.
```

Balanced

```
bool isBalanced(string str) {  
    string ops = operatorsOnly(str);  
    return checkOperators(ops);  
}
```

- Takes in a string str
- Strips down that string to be *just the operators*
- Returns true/false depending on the output of checkOperators -- which takes in the bracket-only string

Balanced -- Your Task

You will be writing two functions:

- 1) `string operatorsFrom(string str)`
 - a) Takes in a string, returns the string with only the brackets
- 2) `bool operatorsAreMatched(string ops)`
 - a) Takes in a string, returns true/false if it's balanced

The catch? You will be implementing both of these recursively!

Remove everything but the Brackets

You've done this iteratively... let's do it recursively:

Similar to the reverse recursive function you've seen in section/lecture!

- Process the first character of the string and determine (if it exists, and) if it should be kept as part of the outputted string
- Recursively process the **rest** of the string (hint: use substrings!)

Be sure to thoroughly test your function before moving on. You should add at least 2-3 tests of your own to validate the behavior of your operatorsFrom function.

operatorsAreMatched

As you think about implementing this recursive function:

A string consisting of only bracketing characters is balanced if and only if one of the following conditions holds:

- The string is empty.
- The string contains "()", "[]", or "{}" as a substring and the rest of the string is balanced after removing that substring.

Let's go over that intuition!

Let's look at the string:

“`[(){}]`”

Let's go over that intuition!

Let's look at the string:

"[(){}]"

We can find the substring "()" inside the above string, so let's remove it



Let's go over that intuition!

Let's look at the string:

“{ }”

Let's go over that intuition!

Let's look at the string:

"[{}]"

We can find the substring "{}" inside the above string, so let's remove it



Let's go over that intuition!

Let's look at the string:

`[]`

Let's go over that intuition!

Let's look at the string:

"[]"



We can find the substring "[]" inside the above string, so let's remove it

Let's go over that intuition!

Let's look at the string:

'''

Let's go over that intuition!

Let's look at the string:

""

We've reached the empty string!! Hmmm... what could this mean??

What happens if we didn't reach the empty string, and couldn't find "(" or "{" or "[" as substrings?? Hmmm...

Questions about Balanced?



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MergeSort

So far, we've been given ways to sort our data structures:

- `v.sort()` for a Vector
- Maps sort their keys automatically in increasing order

Now, it's time for us to take matters into our own hands and use the power of recursion to write a sorting algorithm called **merge sort!**

Binary Merge

Implementing a *binary merge* operation is very common in CS

- Merging two individually sorted sequences into one sorted sequence

You will be implementing:

```
Queue<int> merge (Queue<int> one, Queue<int> two)
```

Binary Merge

```
Queue<int> merge (Queue<int> one, Queue<int> two)
```

- Takes in 2 queues (both in sorted increasing order), and combines them!
 - If they aren't in sorted order throw an error!
- Both queues inputted by value (copy), so feel free to modify them, doesn't matter what they look like after the function is run
- Not necessarily the same lengths
- No special case with duplicates
- To be done **iteratively**, not with recursion.

Binary Merge Tips/Tricks

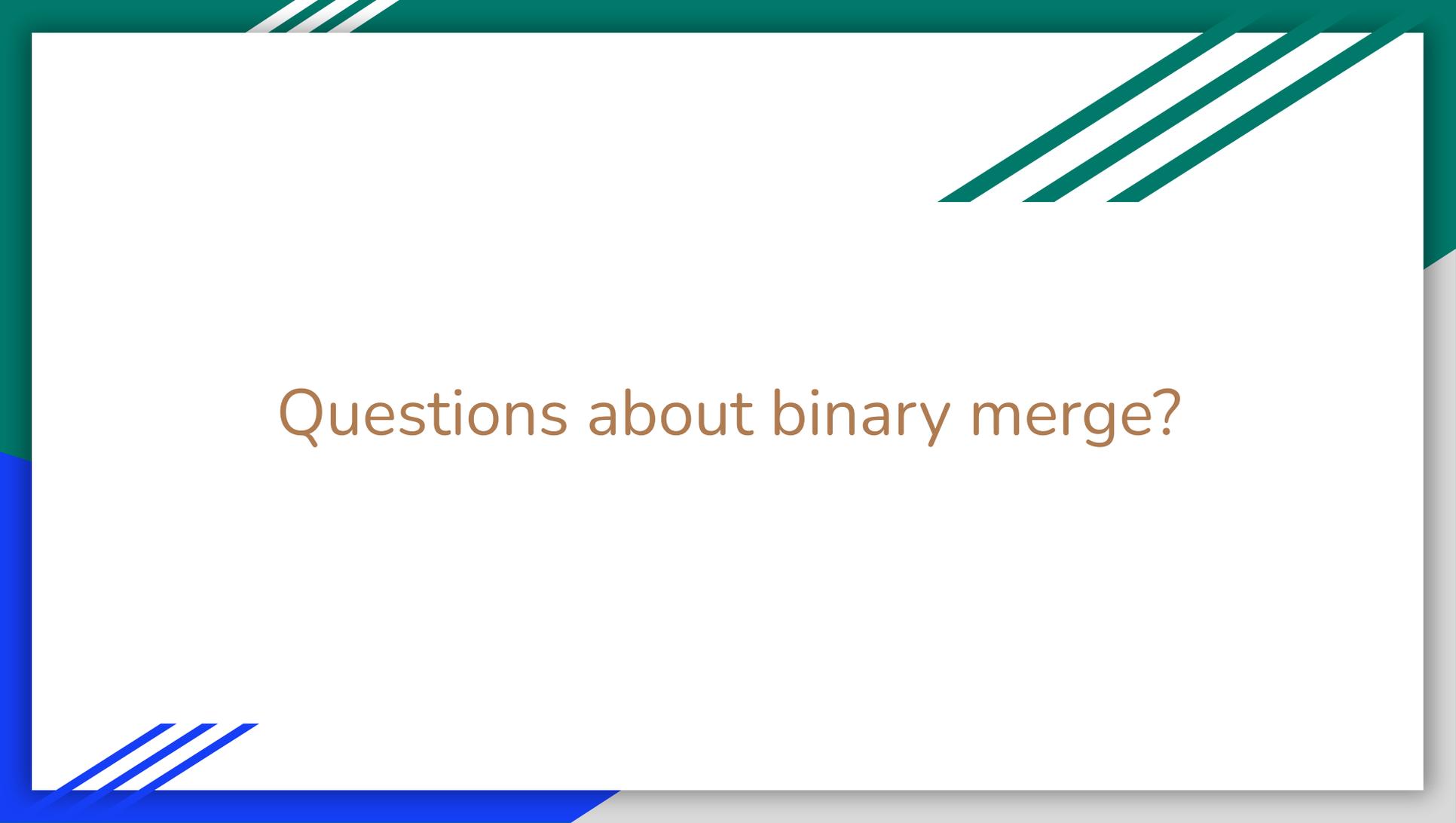
- Go through each queue (iterating through queues is funky remember!)
- Examine the element at the front of each one, enqueue the smaller one into a new Queue (the one you are returning)
- Don't forget about `q.peek()` vs. `q.dequeue()`! Use peek to compare!
- You can either check that queues are in sorted order by:
 - Making a helper that throws an error if it finds an element out of order
 - Do it in your merge function, keeping track of a curr and prev elements!

Big-O Analysis of Iterative Merge

You will be asked to predict/calculate a Big-O analysis of the Iterative Merge method!

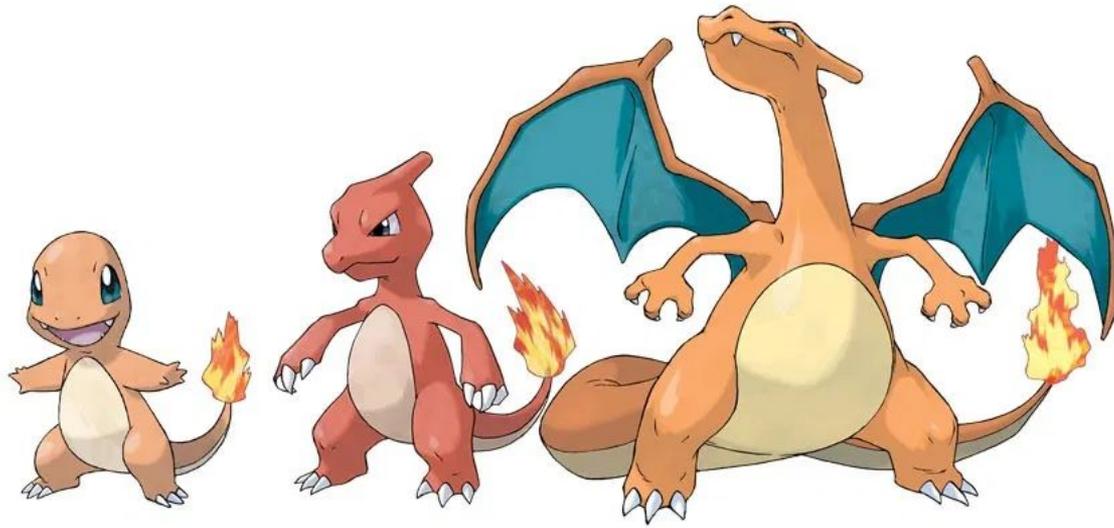
- Use Time Trials here! Find the relationship between size of queues and times!
- Experiment with varying sizes

Q11. Include the data from your execution timing and explain how it supports your Big O prediction for binary merge.



Questions about binary merge?

We can do better than that! Level up!



We can do better than that! Level up!

RECURSION



Recursive MultiMerge

A multi-way merge takes the "divide-and-conquer" strategy to recursively merge a sequence.

Especially for longer sequences, this approach is **much more efficient**.

Our job is to write

```
Queue<int> recMultiMerge(Vector<Queue<int>>& all)
```

Recursive MultiMerge Algorithm

The handout says it best:

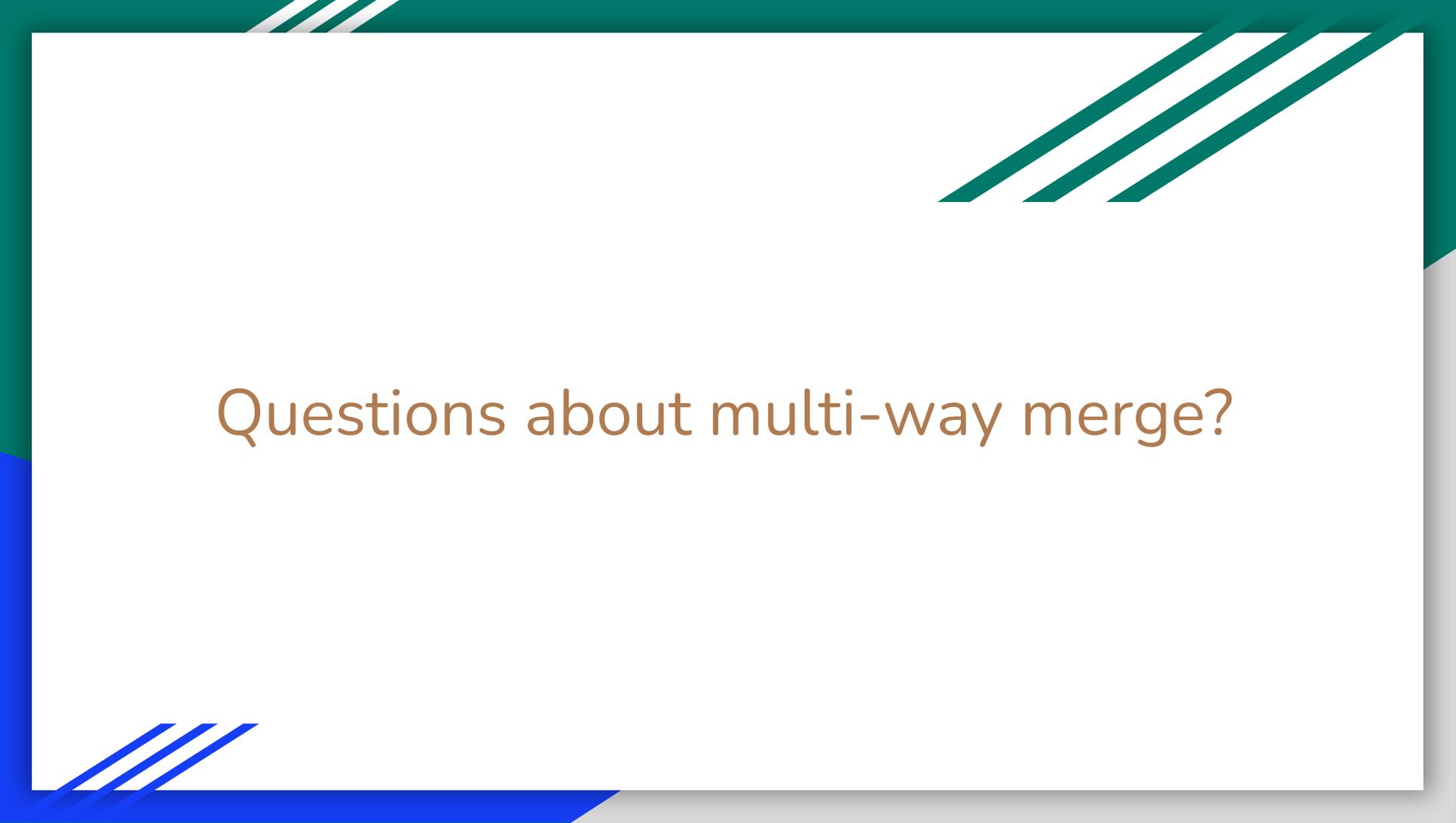
1. Divide the input collection of k sequences into two halves. The "left" half is the first $K/2$ sequences in the vector, and the "right" half is the rest of the sequences.
 - The `Vector subList` operation can be used to subdivide a Vector, which you may find helpful.
2. Make a recursive call to `recMultiMerge` on the "left" half of the sequences to generate one combined, sorted sequence. Then, do the same for the "right" half of the sequences, generating a second combined, sorted sequence.
3. Use your binary `merge` function to join the two combined sequences into the final result sequence, which is then returned.

Implementation Tips

- Consider your base cases:
- You take in a `Vector<Queue<int>>` as a parameter, right?
 - What are the smallest possible cases we might get for an input of this type?
 - How should we deal with them?

Implementation Tips

- Trust the recursion.
 - Recursively merge the left and the right
 - Merge the 2 final resulting queues into 1!
- Remember: your recursion should lead you to your base cases
- Recursion is all about breaking down our problem into smaller pieces
 - What are the pieces that are getting smaller with each recursive call?
- As long as you follow the handout, you should be in good shape!



Questions about multi-way merge?

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A meme featuring Aragorn from The Lord of the Rings. He has a serious expression and is looking slightly to the side. A large red 'X' is drawn over his mouth, indicating that the text is being corrected or negated. The background is a warm, golden-brown color, likely from a scene in the movie.

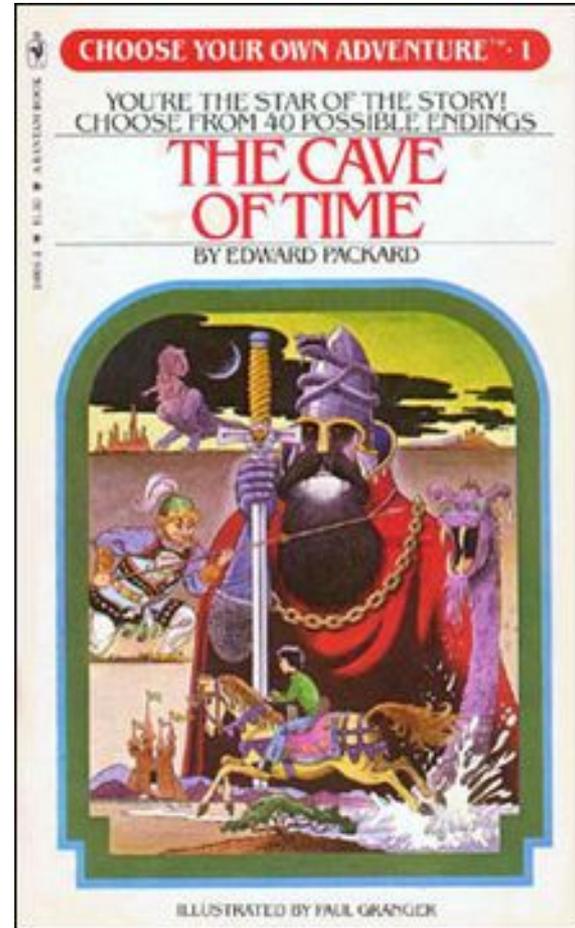
One does not
simply...

bactrack...

Remember these? →

Choose-Your-Own-Adventure can be thought of as recursive backtracking.

We **choose** one path, **explore** it, then **unchoose** it if we do not reach a solution/happy ending.



Recursive Backtracking

Recursive backtracking helps us solve problems for which we need to try many different options to find the desired solution.

If we make a step in the wrong direction, we need to backtrack -- undo that step -- and choose a different step.

To construct a backtracking solution, you'll likely have to recurse on all possible options. If you reach the end of an option, you've found a solution!





The Game of Boggle

In this game, each player must try to find as many words as possible on the Boggle Board. Each word has a corresponding score depending on its length.

We will be coding a computerized Boggle player, whose goal is to find the maximum possible score!

The Game of Boggle

Of course, we will harness the power of recursive backtracking!

- Look at each letter on the board
- Find all words that begin with that letter

```
int scoreBoard(Grid<char>& board, Lexicon& lex)
```

The function operates by recursively exploring the board to find all words and tallying up the points.

Boggle Rules

- A word can only build up using a character that is adjacent (4 cardinal directions NSEW, **and** diagonals NE, SE, NW, SW)
- Each word has a score that contributes to the point total



Big Idea 1: Backtracking

Our backtracking solution should recurse on **all possible options**, since we need to examine every possible word on the board.

This means every single letter-cube on the board (every single location on the Grid) must be examined as the first letter of a word.

Then, we can proceed in the adjacent cubes around that letter-cube and build up words!

Scoring

- A word must be at least 4 letters long to be scored
 - length 4 → 1 point
 - length 5 → 2 points
 - length 6 → 3 points ... etc.
- Only *unique* words should contribute to the score
 - Even if the word "hello" happens to appear in 3 different variations on the board, the word only counts once in the score total

Can this word-so-far contribute to my score?

Yes, if...

- It is at least 4 letters long
- It is a valid word, according to a function from the [Lexicon class](#)
- The word can be traced through adjacent letter cubes
- The word has not yet been seen / contributed to the score total

Big Idea: Marking

By this point, hopefully it's clear that we should be keeping track of the words we've already scored.

Hint: we will also need to keep track of locations we've already visited on the board, to avoid going back to a character that has already been added to the word.

This means we'll be "marking" locations as seen as we reach them.

Once we've reached a word...

- Great!
- Food for thought: should we stop recursing here?
 - Example: if I've found "TREAT" and I stop here,
 - will I ever find "TREATY"?



Some tips to keep in mind...

- Using **GridLocations** makes it much easier to look at any given point on the board (a Grid)
- Keep track of visited locations using a data structure that doesn't allow duplicates 🤔
- Careful not to go out of bounds... **g.inBounds()** sounds nice here 🤔
- Take advantage of the functions in the Lexicon class that give you information about a specified word!
 - Very helpful in determining full words and prefixes so we can **prune**

Prune, you say?

10 Health Benefits of... Prunes

1. Assists Iron Absorption
2. Good Vitamin C Source
3. Normalise Blood Sugar
4. Prevents Over-eating
5. Lowers Cholesterol
6. Improve Bone Health
7. Protects Intestines
8. Disease Protection
9. Natural Laxative
10. Anti-Oxidant



EatHealthyLiveFit.com



Prune, you say?

Backtracking

Benefits of...

Prunes

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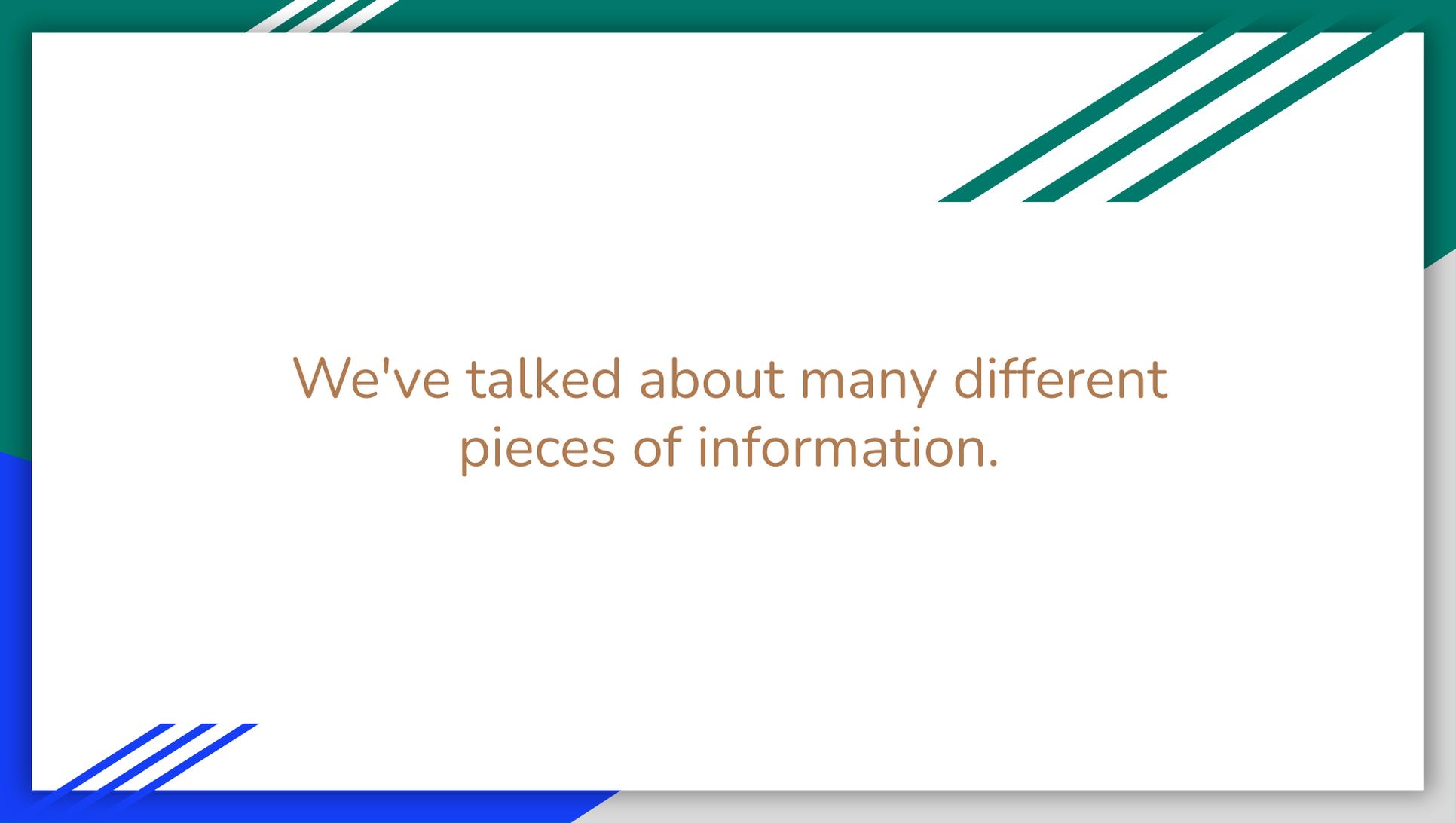
Pruning

- Say you are building up a word, and that word so far is "xbqmah".
- Even if there are still more letter blocks to visit on the Boggle board, *what's the point?*
 - We'd be doing ourselves (and our program) a disservice by continuing down a path that we know is a dead end from the start.
- The good news is, the `Lexicon` class has a function **`containsPrefix(wordSoFar)`** that tells us whether some `wordSoFar` is a prefix of any word. Be sure to make use of it!

WAIT A MINUTE

WHO ARE YOU?

FlukyFeed



We've talked about many different pieces of information.

We've talked about many different pieces of information.

But the only information we get is the 2 parameters:

```
int scoreBoard(Grid<char>& board, Lexicon& lex)
```

That isn't enough to keep track of all the information we need!

The helper function - wrapper function relationship is **key** in recursive backtracking.

Often, we need to keep track of information as we go down a path, meaning we need to have all of the appropriate parameters.

Consider how you might use this approach in your solution!

- What pieces of information do I need to have access to within each function call?

An infamous meme from Trip:



What should we call this new word
game?

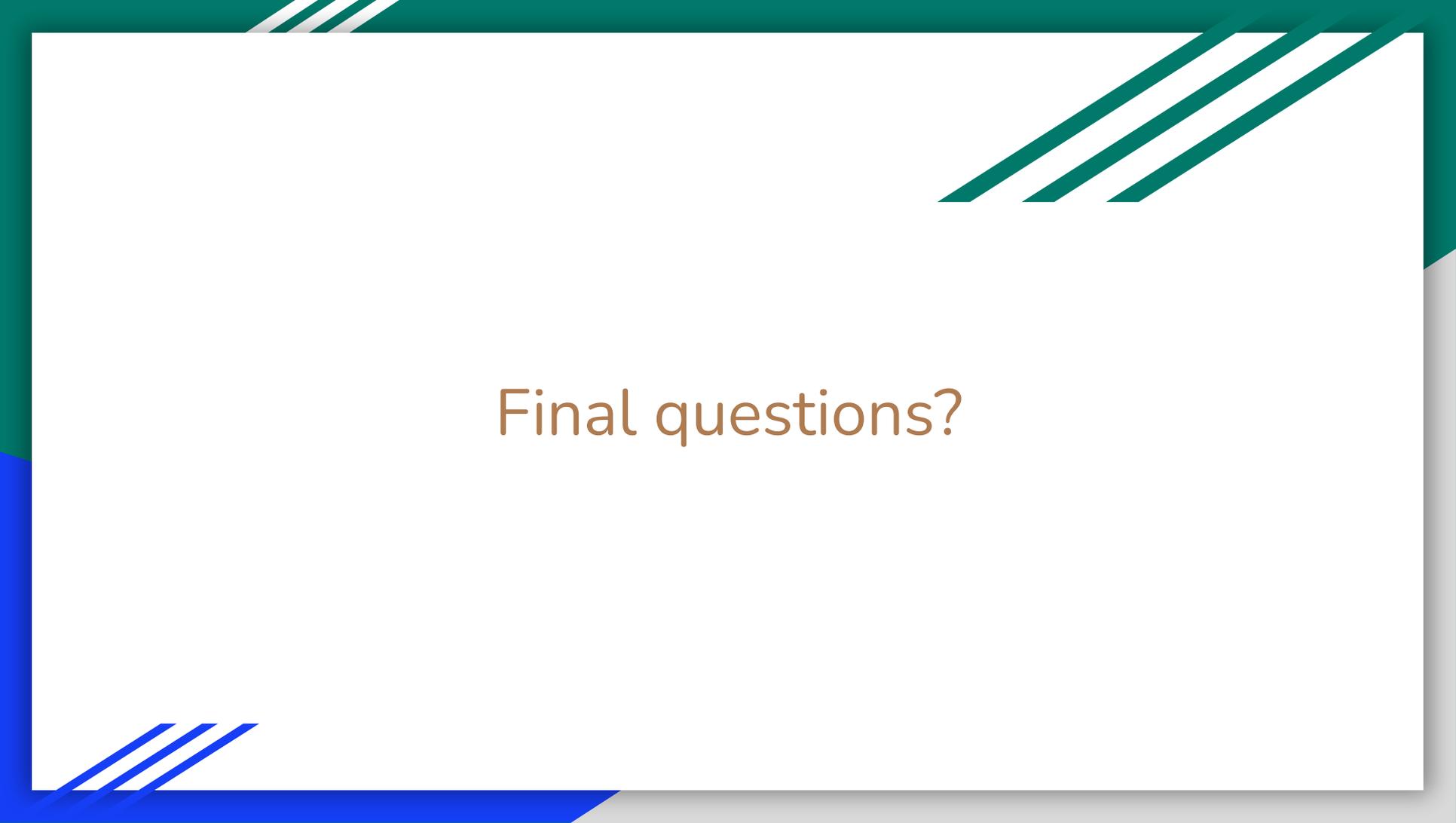
Dog: how about Doggle?

Bog: I have a better idea





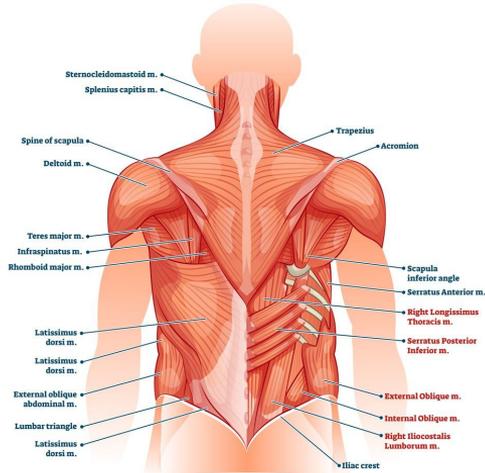
Questions about Boggle?



Final questions?

Good luck

LOWER BACK MUSCLES



-ing!