

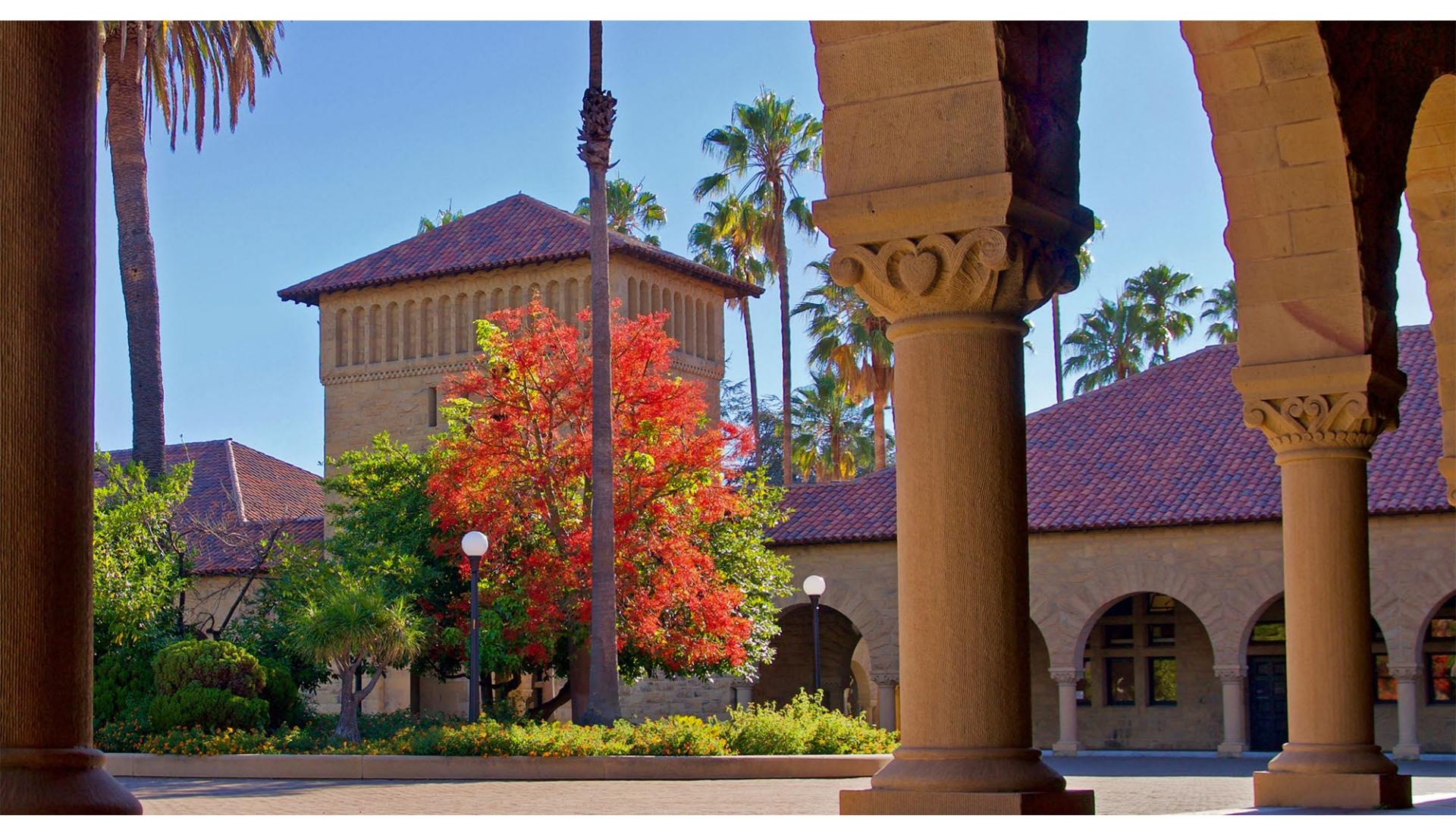


# Binary Search Trees

What is your favorite type of tree?  
(e.g., oak, redwood, maple)

put your answers in the chat











# Roadmap

## Object-Oriented Programming

Roadmap graphic courtesy of Nick Bowditch

### C++ basics

User/client

vectors + grids

stacks + queues

sets + maps

Core Tools

testing

algorithmic analysis

recursive problem-solving

Diagnostic

Life after CS106

real-world algorithms

Implementation

arrays

dynamic memory management

linked data structures

# Roadmap

## Object-Oriented Programming

### C++ basics

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# Today's questions

How can we take  
advantage of trees to  
structure and efficiently  
manipulate data?

# Today's topics

1. What is a binary search tree (BST)?
2. Building efficient BSTs

1. Implementing Sets with BSTs

# Review

[tre e s ]

# Definition

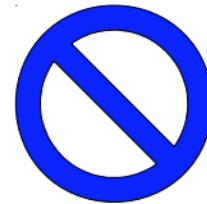
## tree

A tree is hierarchical data organization structure composed of a root value linked to zero or more non-empty subtrees.

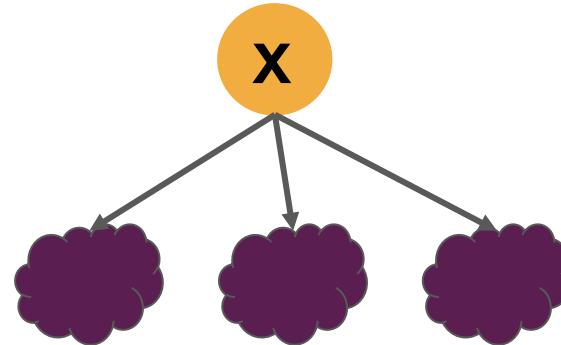
# What is a tree?

A tree is either...

An empty data structure, or...



A single node (parent), with zero or more non-empty subtrees (children)

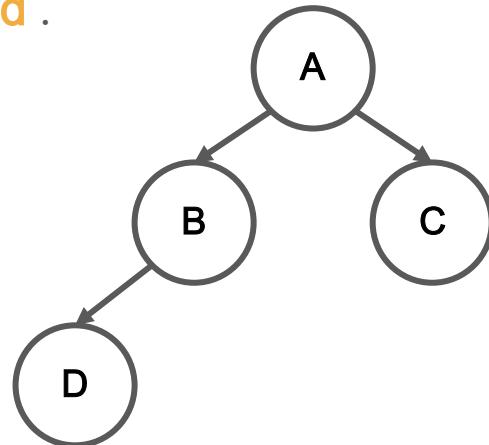


# Tree terminology

- Types of nodes
  - The **root** node defines the "top" of the tree.
  - Every node has 0 or more **children** nodes descended from it.
  - Nodes with no children are called **leaf nodes**.
  - Every node in a tree has exactly one **parent** node (except for the root node).
- Terminology for quantifying trees
  - A **path** *between two nodes* traverses edges between parents and their children, and **length of a path** is the number of edges between the two nodes.
  - The **depth of a node** is the length of the path (# of edges) between the root and that node.
  - The **height of a tree** is the number of nodes in the longest path through the tree (i.e. the number of **levels** in the tree).

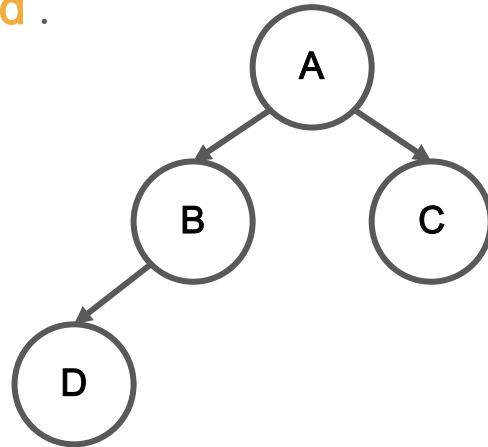
# Binary trees

- A **binary tree** is a tree where every node has either 0, 1, or 2 children. No node in a binary tree can have more than 2 children.
- Typically, the two children of a node in a binary tree are referred to as the **left child** and the **right child** .



# Binary trees

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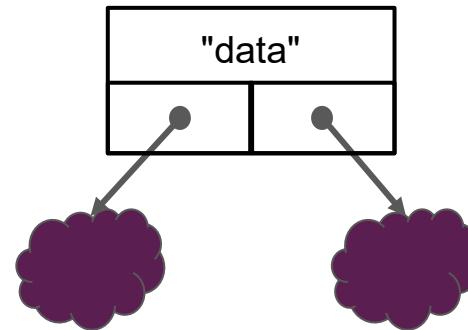
```
struct TreeNode {  
    string data;  
    TreeNode* left;  
    TreeNode* right;  
}
```

# What is a tree in C++?

A tree is either...

An empty tree  
represented by  
`nullptr`, or...

A single `TreeNode`,  
with 0, 1, or 2 non-  
null pointers to  
other `TreeNodes`



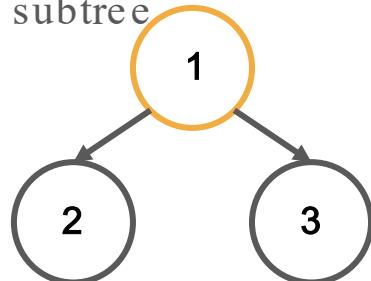
# Building a tree

- Building a tree is very similar to the process of building a linked list.
- We create new nodes of the tree by dynamically allocating memory.
- We start by first creating the leaf nodes and then creating their parents.
- We integrate the parents into the tree by rewiring their **left** and **right** pointers to the already-created children.

# Traversing a tree - recursively!

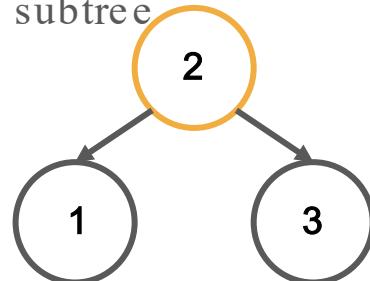
## Pre-order

1. **"Do something" with the current node**
2. Traverse the left subtree
3. Traverse the right subtree



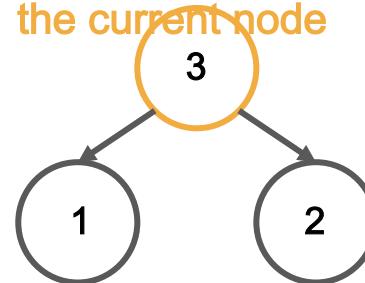
## In-order

1. Traverse the left subtree
2. **"Do something" with the current node**
3. Traverse the right subtree



## Post-order

1. Traverse the left subtree
2. Traverse the right subtree
3. **"Do something" with the current node**



# Freeing a tree!

1. What kind of traversal would you use?
2. Where does the delete call go?

# Let's code it!

**freeTree()**

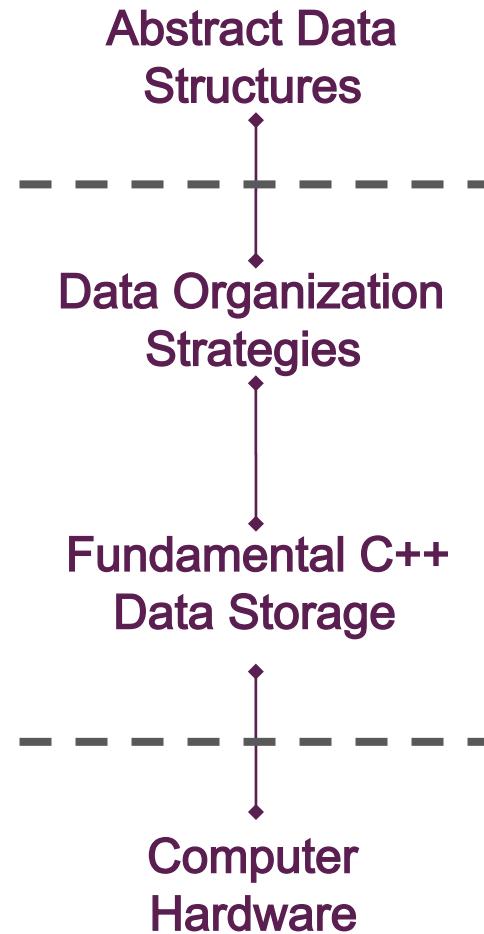
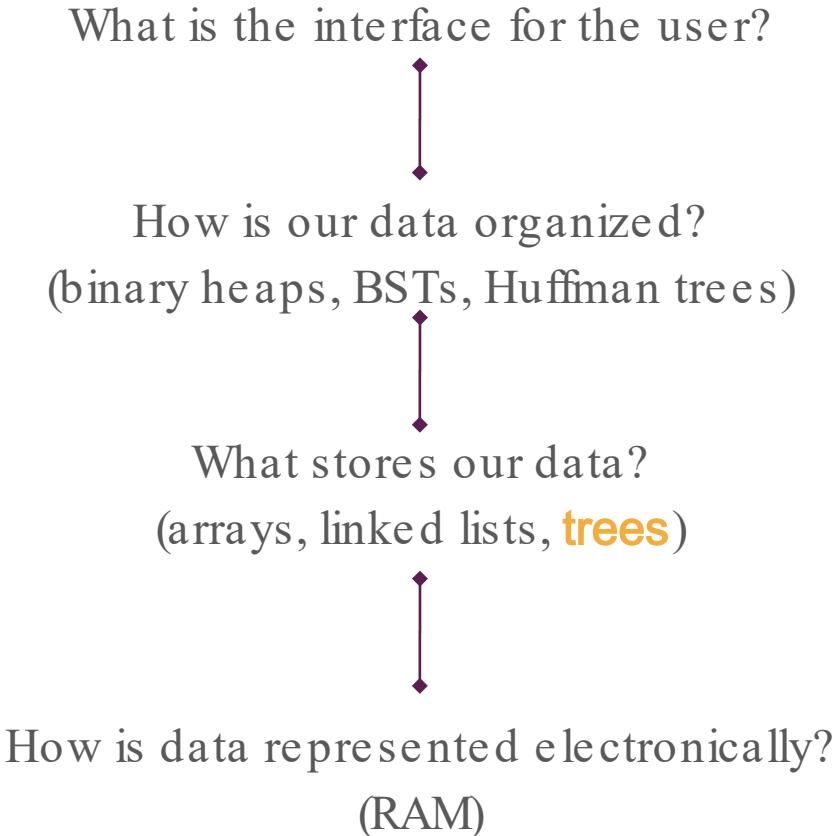
**Key Idea:** The distance from each element (node) in a tree to the top of the tree (the root) is small, even if there are many elements.

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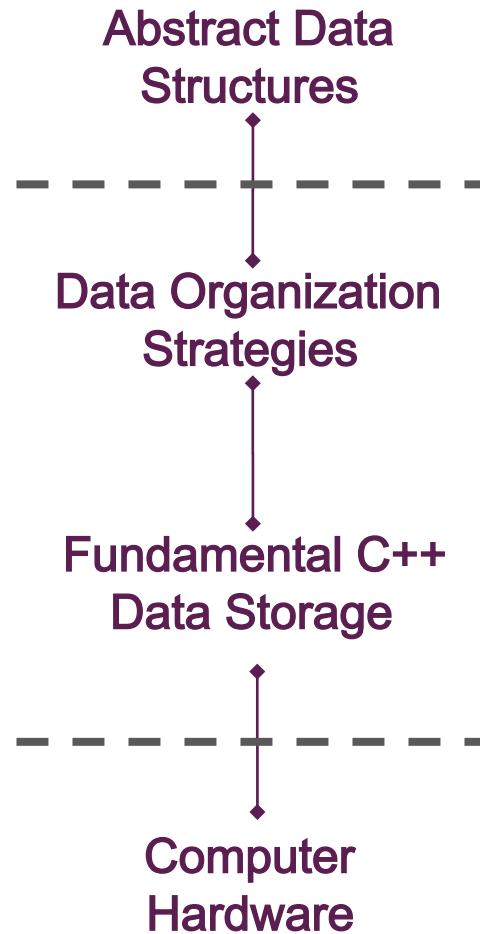
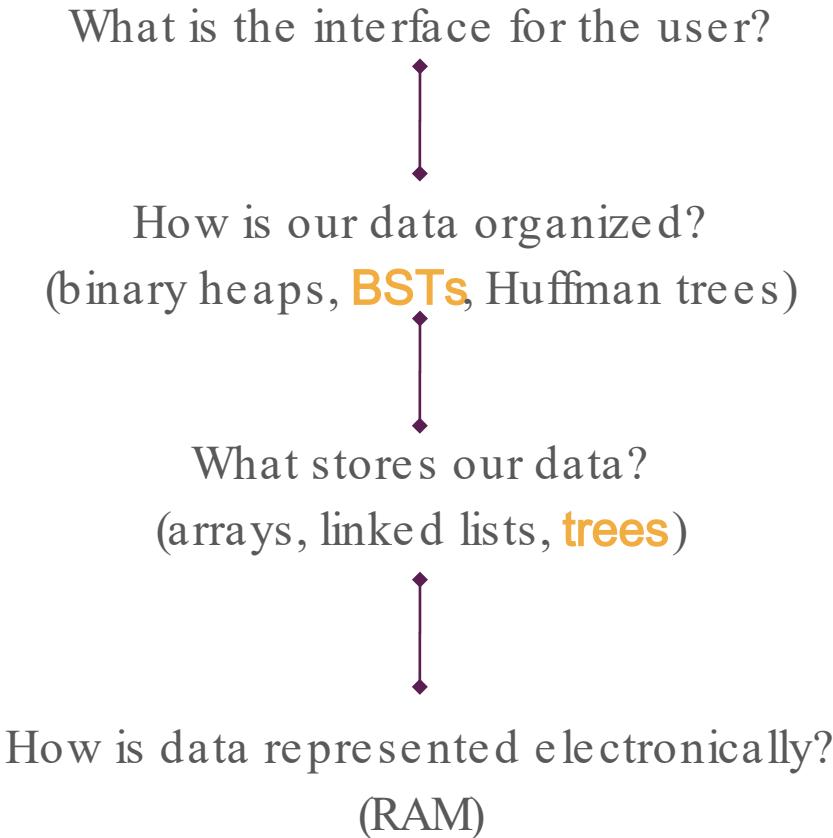
How can we take advantage of trees to structure and manipulate data?

Revisiting our levels of abstraction...

# Levels of abstraction



# Levels of abstraction



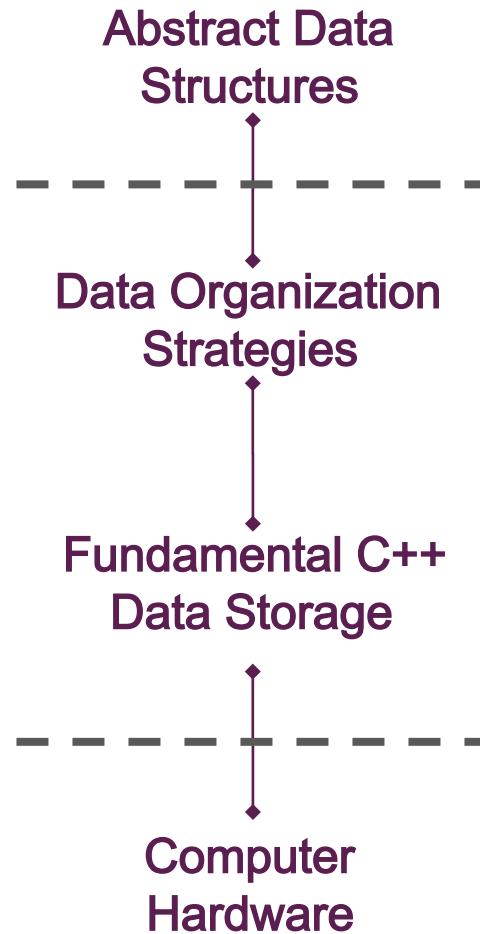
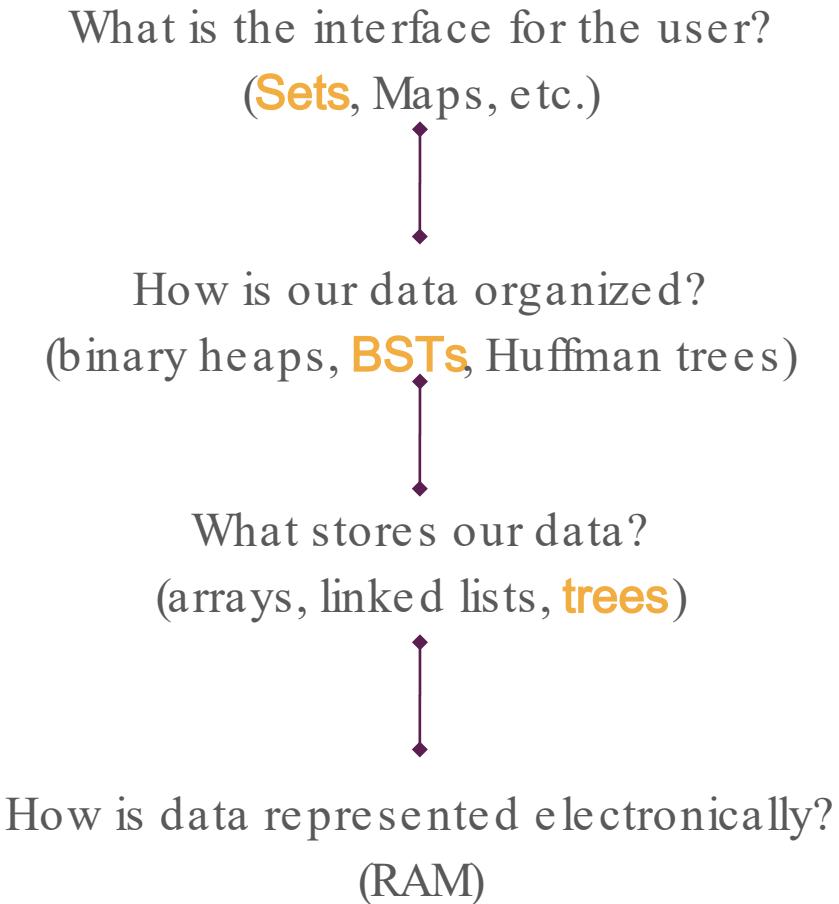
# ADT Big-O Matrix

- Vectors
  - `.size()` -  $O(1)$
  - `.add()` -  $O(1)$
  - `v[i]` -  $O(1)$
  - `.insert()` -  $O(n)$
  - `.remove()` -  $O(n)$
  - `.clear()` -  $O(n)$
  - `traversal` -  $O(n)$
- Grids
  - `.numRows() / .numCols()` -  $O(1)$
  - `g[i][j]` -  $O(1)$
  - `.inBounds()` -  $O(1)$
  - `traversal` -  $O(n^2)$
- Queues
  - `.size()` -  $O(1)$
  - `.peek()` -  $O(1)$
  - `.enqueue()` -  $O(1)$
  - `.dequeue()` -  $O(1)$
  - `.isEmpty()` -  $O(1)$
  - `traversal` -  $O(n)$
- Stacks
  - `.size()` -  $O(1)$
  - `.peek()` -  $O(1)$
  - `.push()` -  $O(1)$
  - `.pop()` -  $O(1)$
  - `.isEmpty()` -  $O(1)$
  - `traversal` -  $O(n)$
- Sets
  - `.size()` -  $O(1)$
  - `.isEmpty()` -  $O(1)$
  - `.add()` -  $O(\log(n))$
  - `.remove()` -  $O(\log(n))$
  - `.contains()` -  $O(\log(n))$
  - `traversal` -  $O(n)$
- Maps
  - `.size()` -  $O(1)$
  - `.isEmpty()` -  $O(1)$
  - `m[key]` -  $O(\log(n))$
  - `.contains()` -  $O(\log(n))$
  - `traversal` -  $O(n)$

# ADT Big-O Matrix

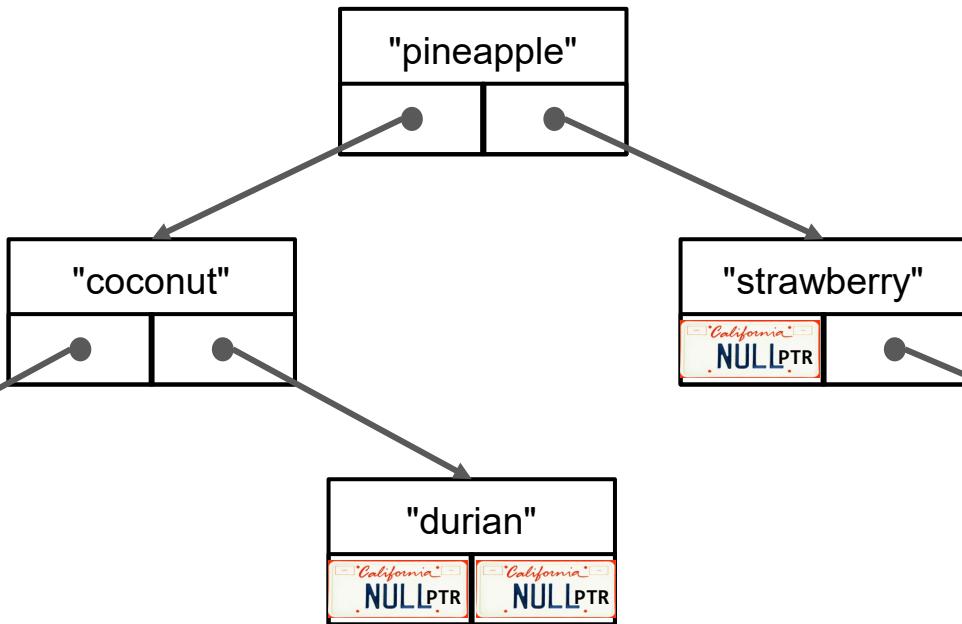
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## Levels of abstraction



What is a binary search tree  
(BST)?

# Building Trees Programmatically

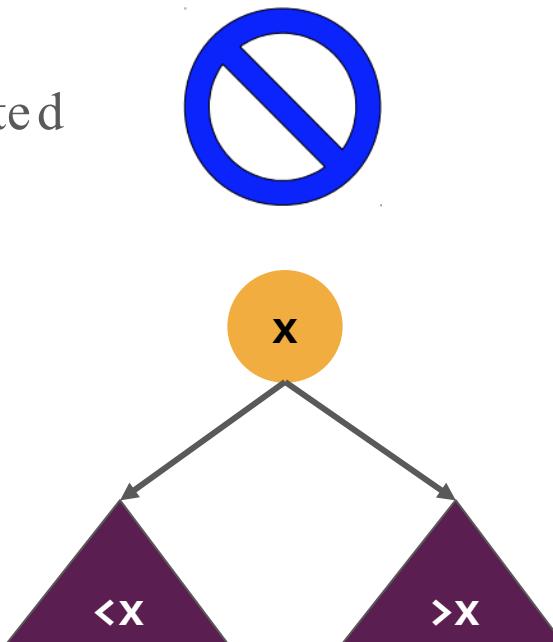


```
struct TreeNode {  
    string data;  
    TreeNode* left;  
    TreeNode* right;  
}
```

# A binary search tree is either...

an empty data structure represented by `nullptr` or...

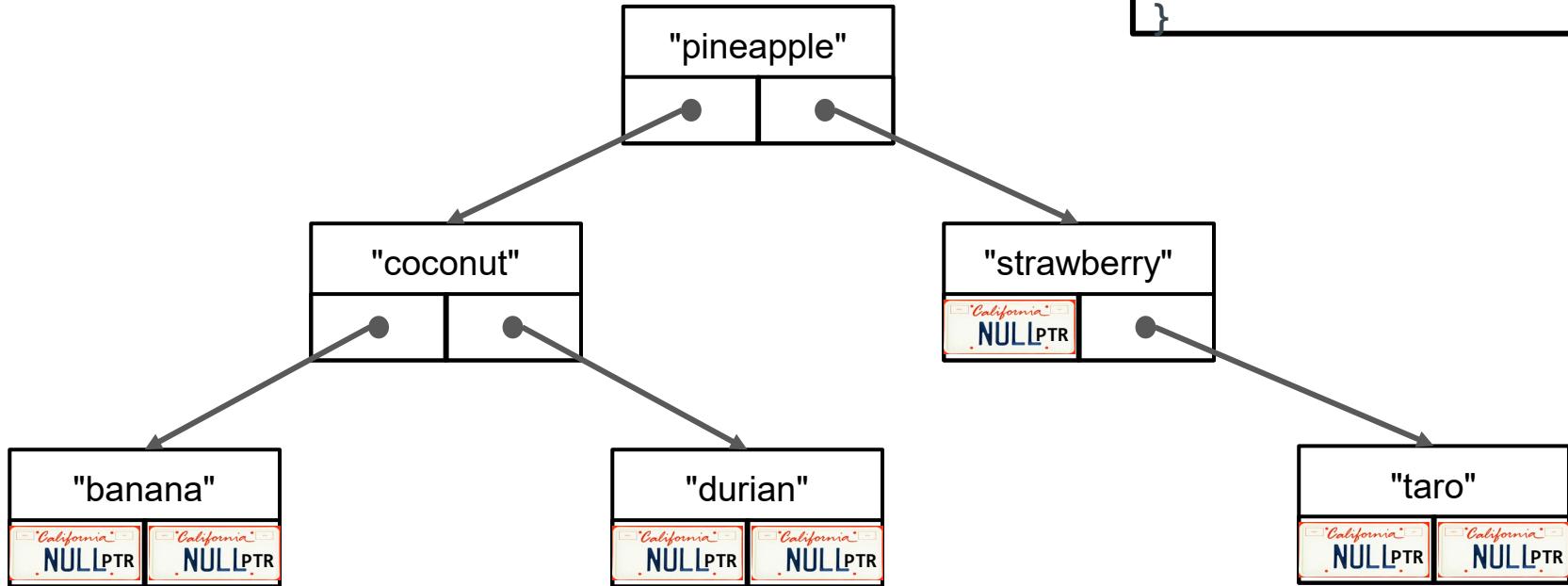
a single node, whose left subtree is a BST of smaller values than **x**...



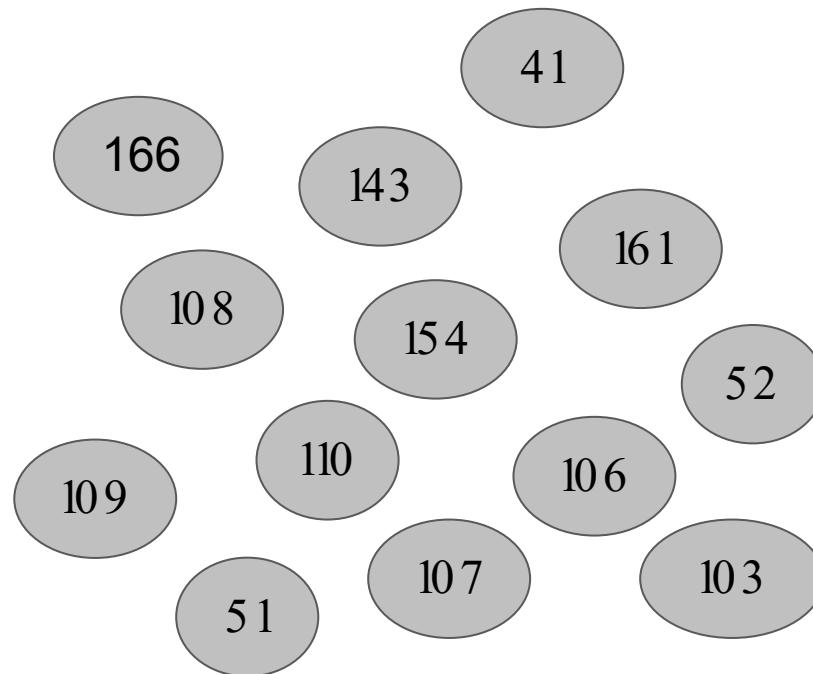
and whose right subtree is a BST of larger values than **x**.

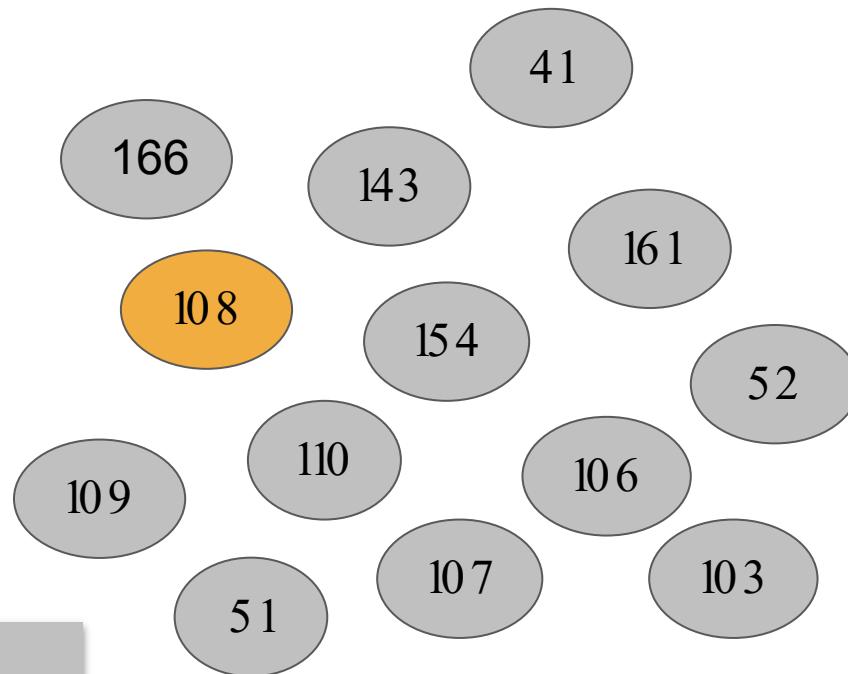
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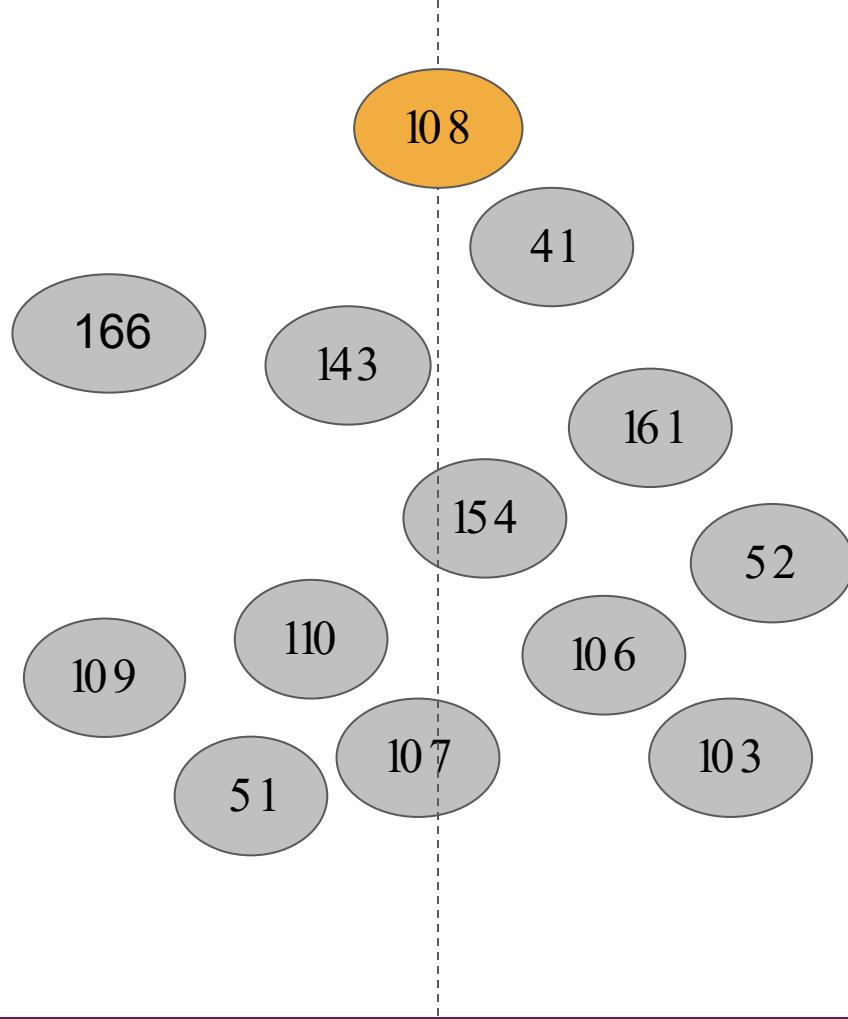


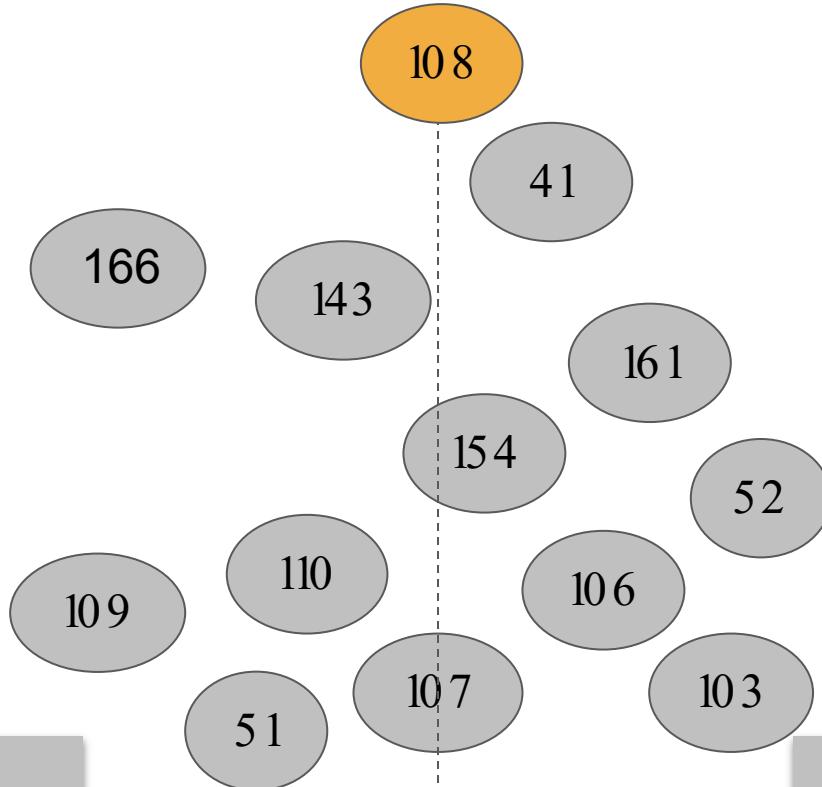
# Building a BST

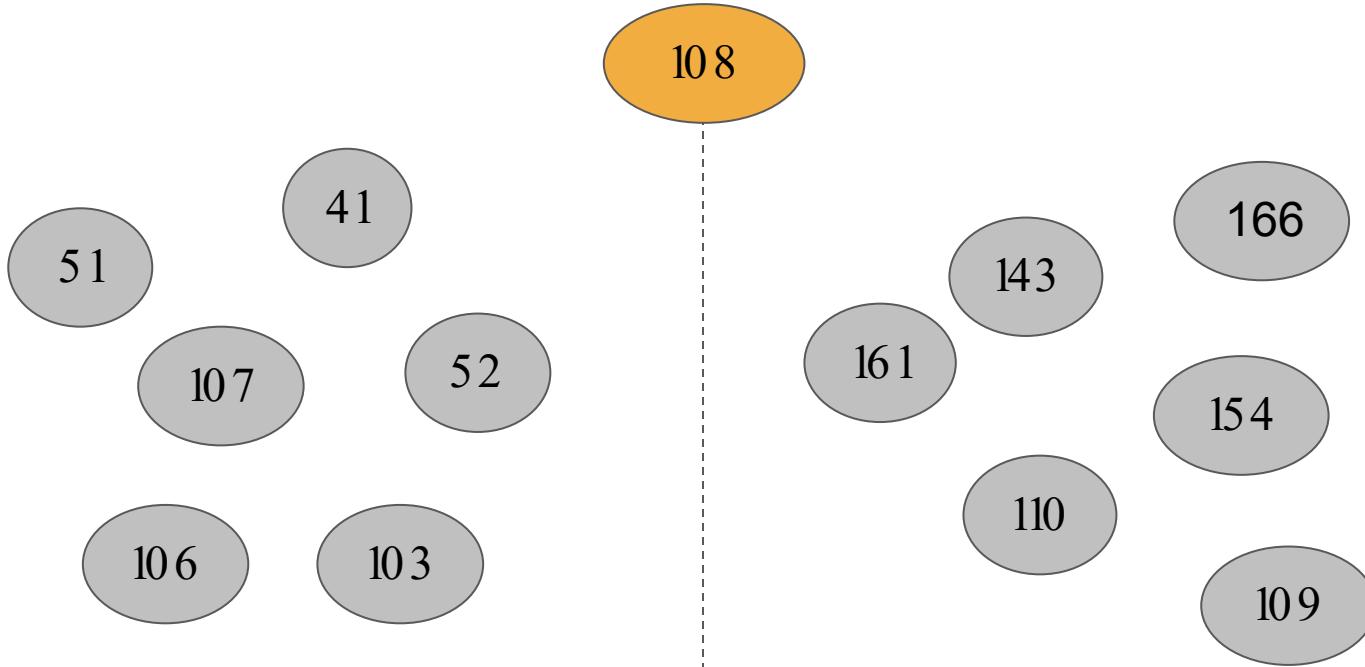




Pick the median element

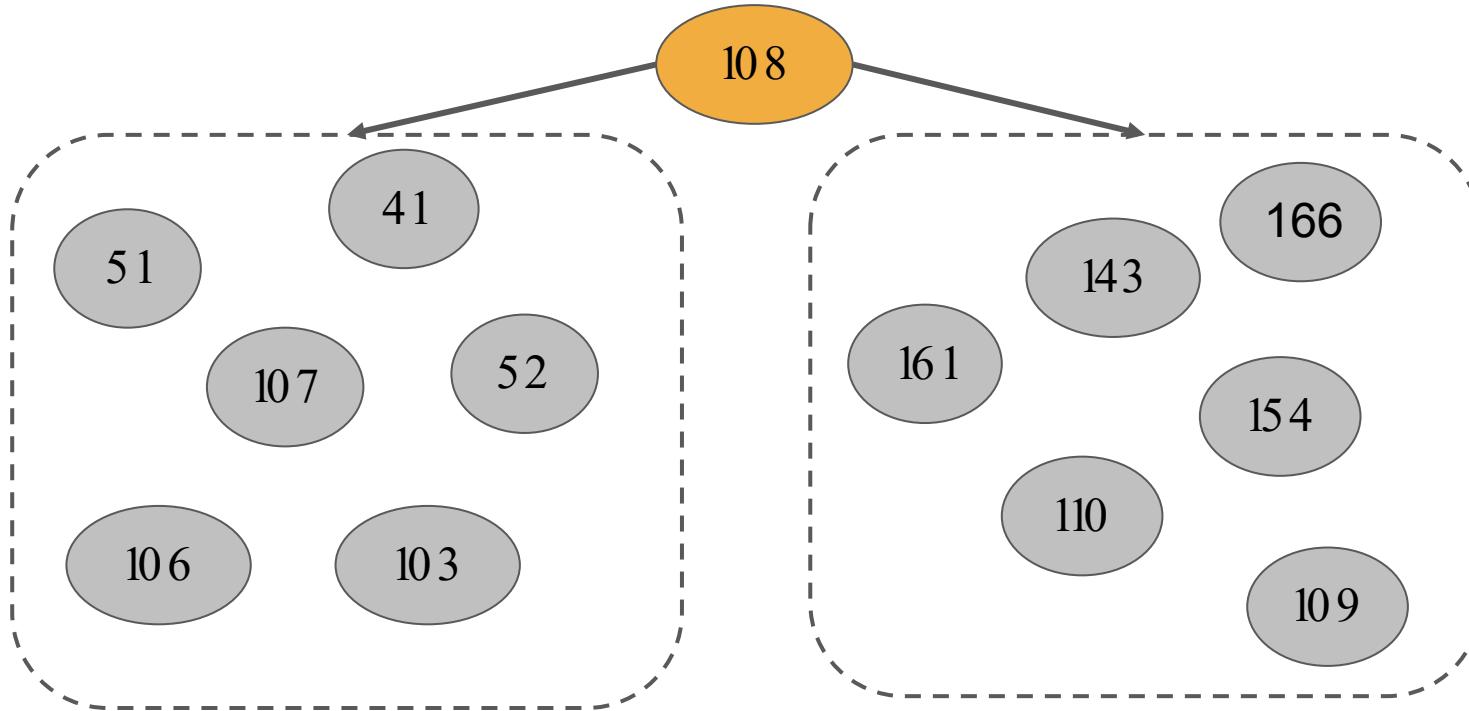


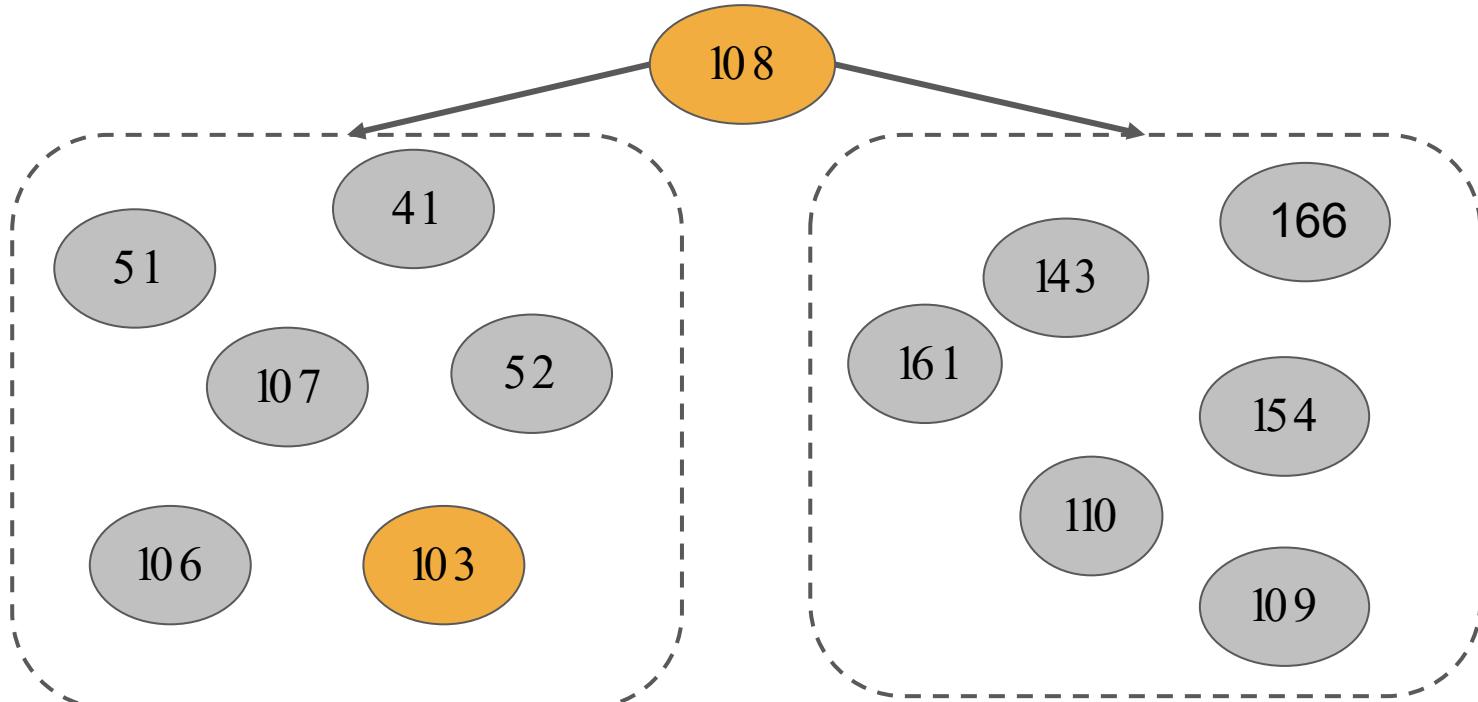




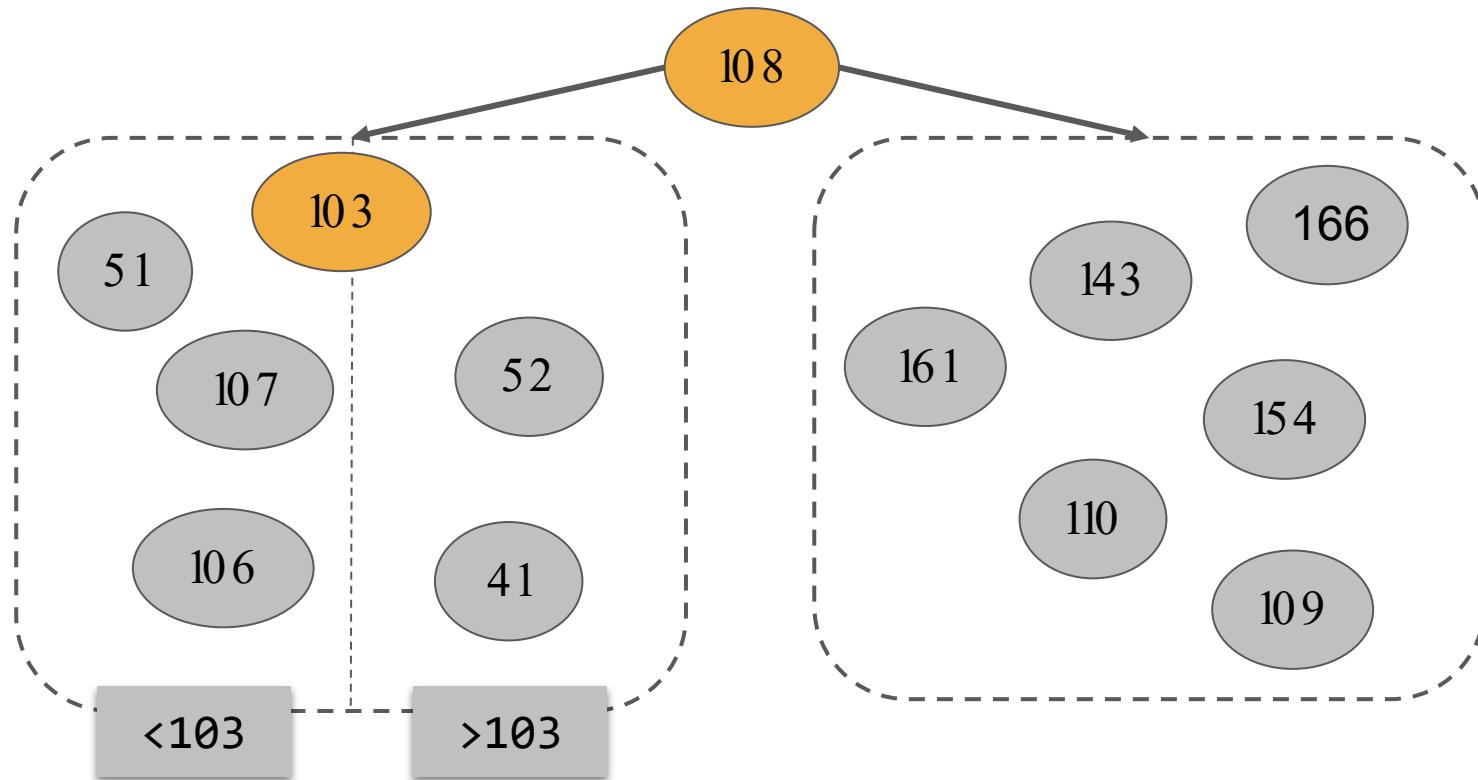
Move elements less than 10 8 to this side

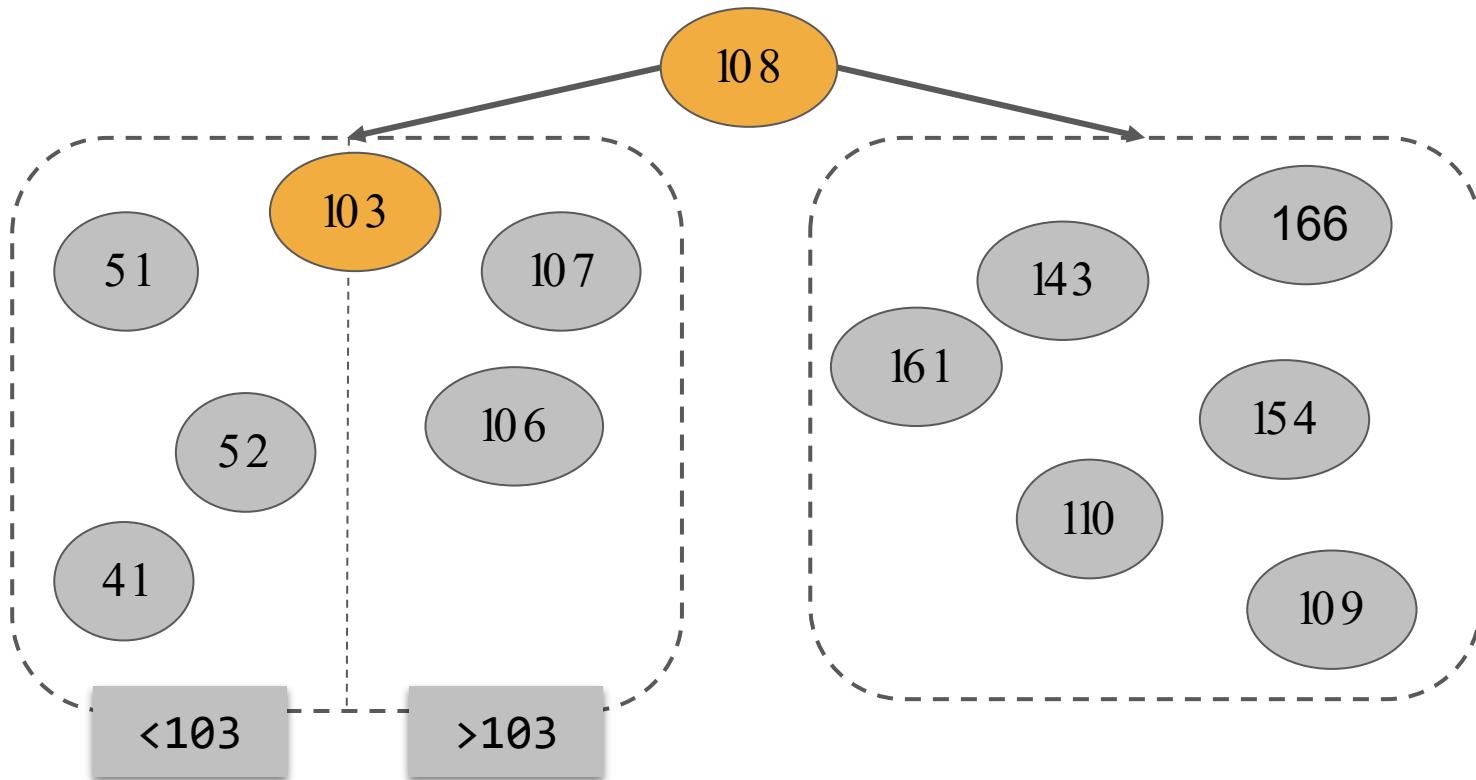
Move elements greater than 10 8 to this side

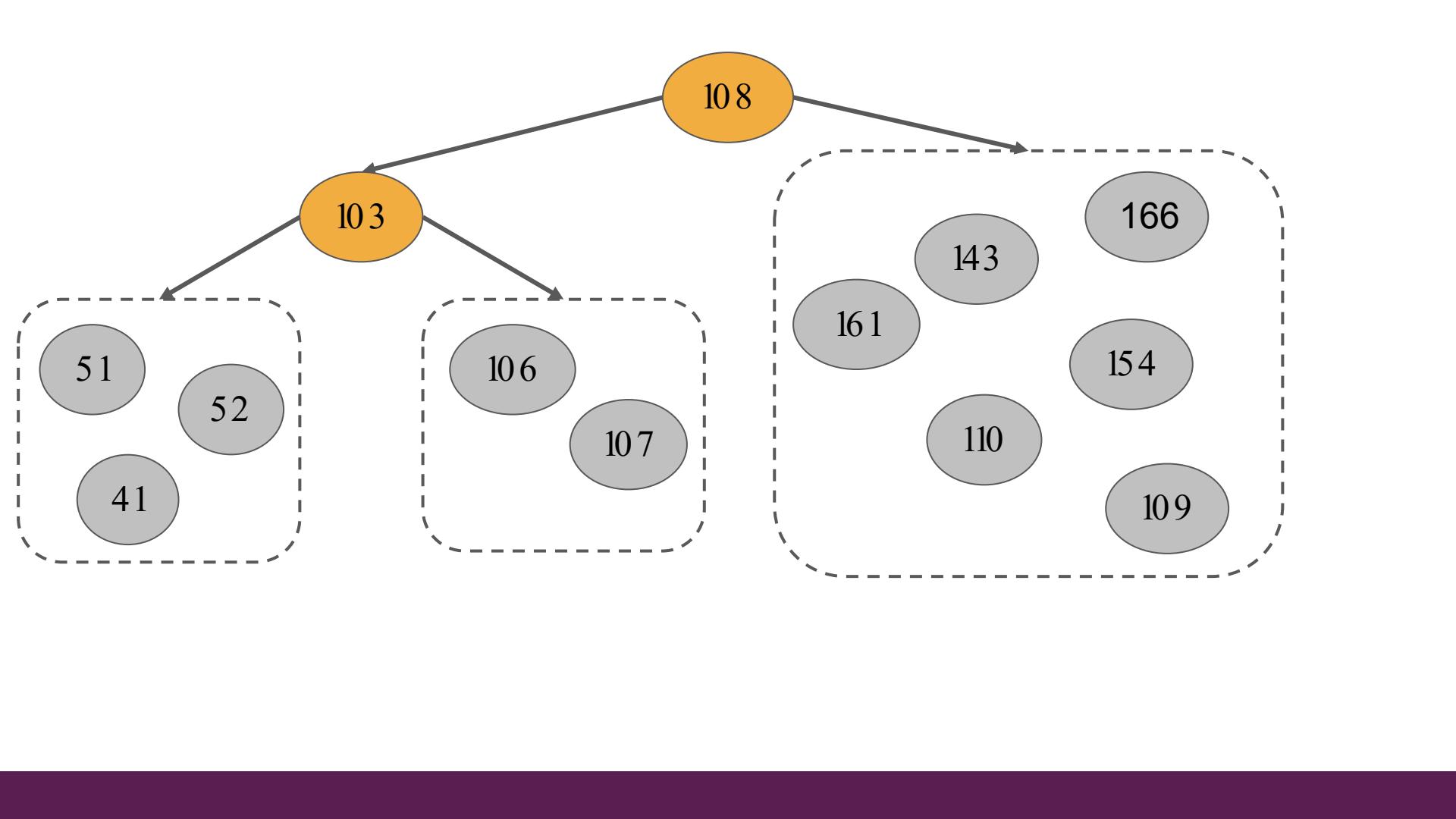


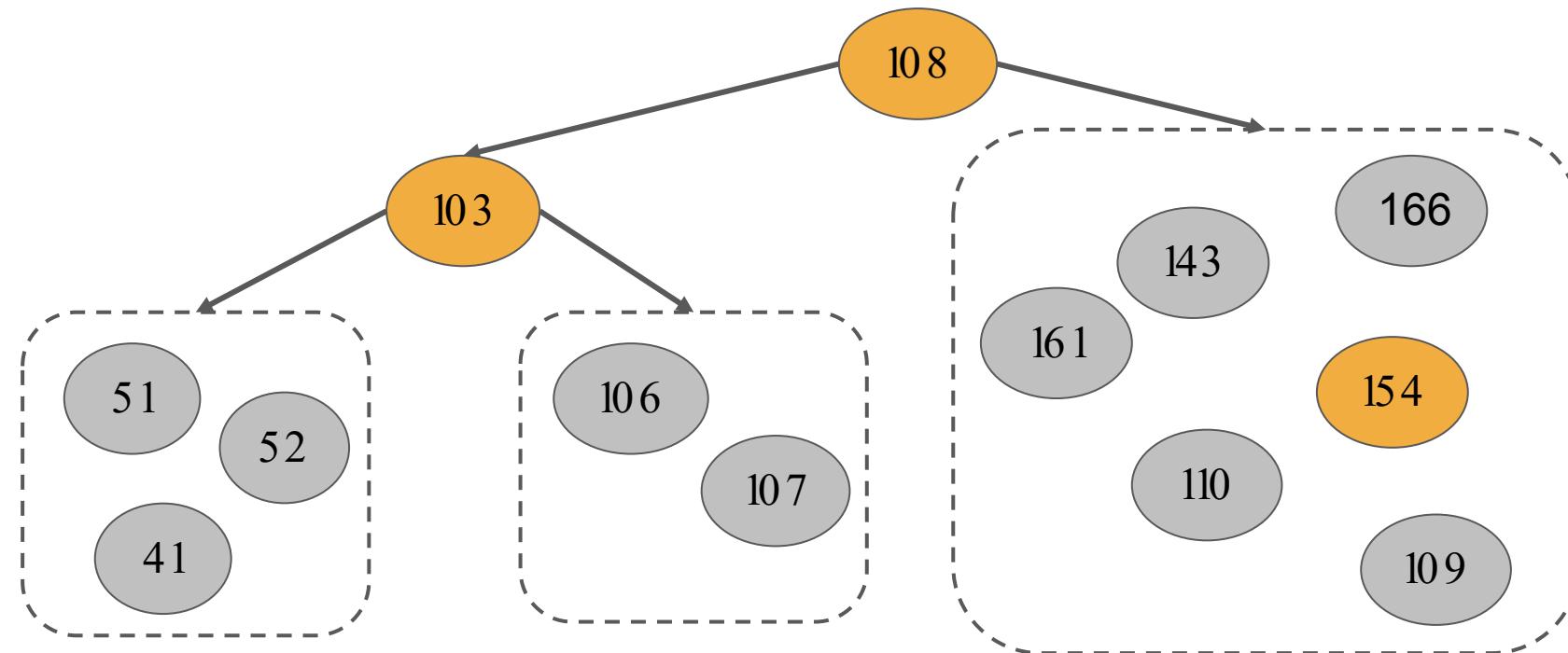


Pick the median element  
of the left side

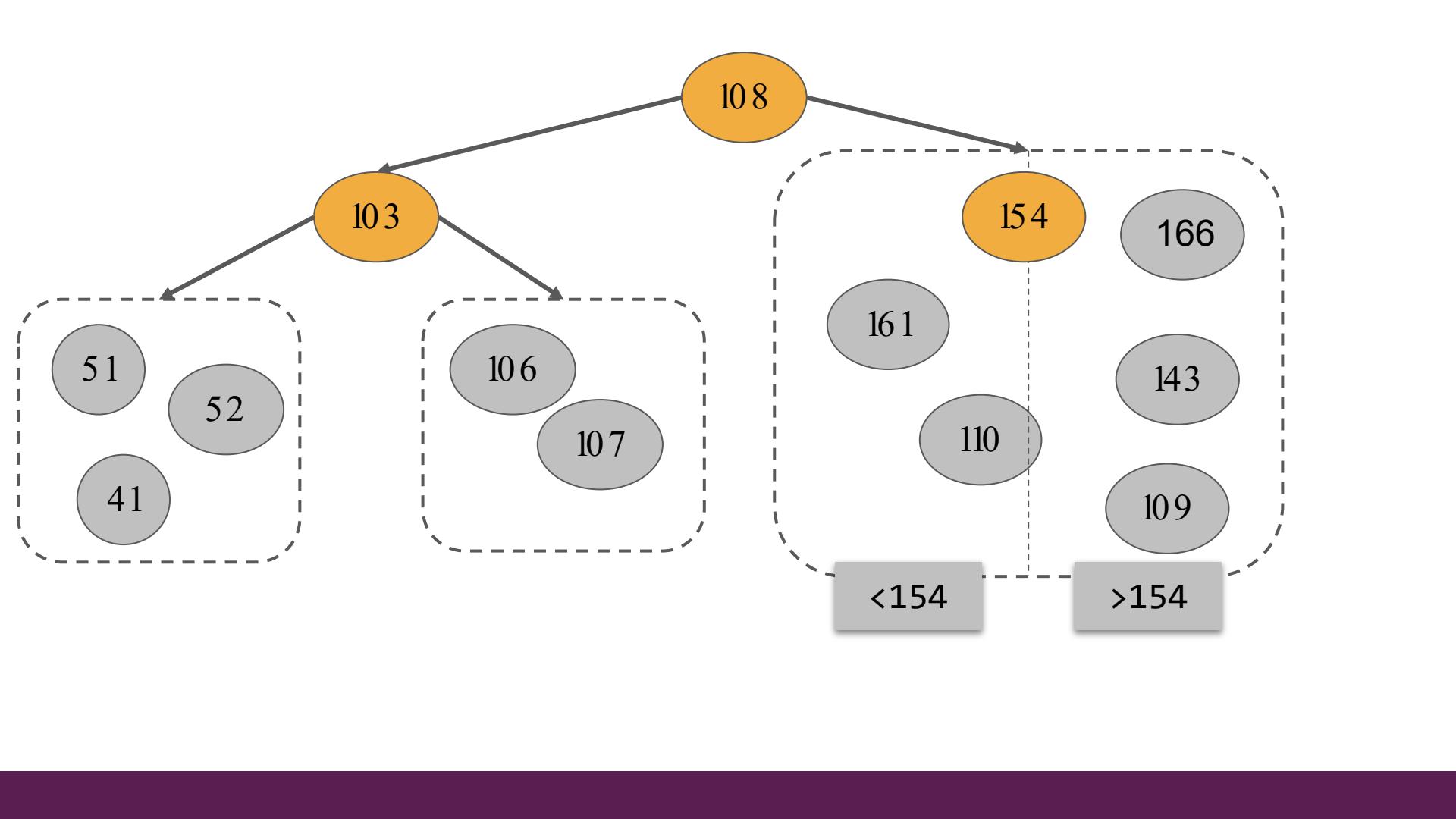


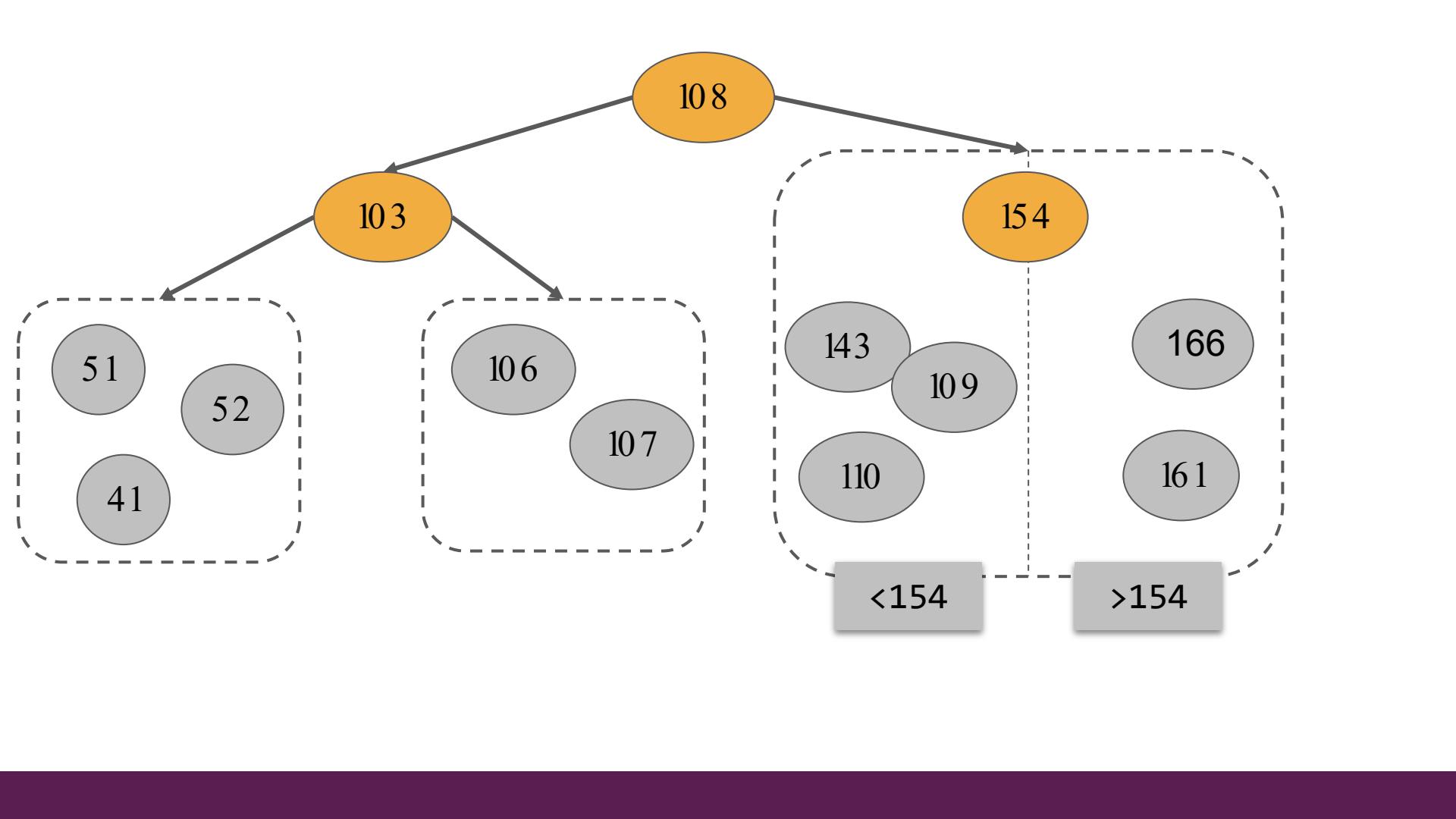


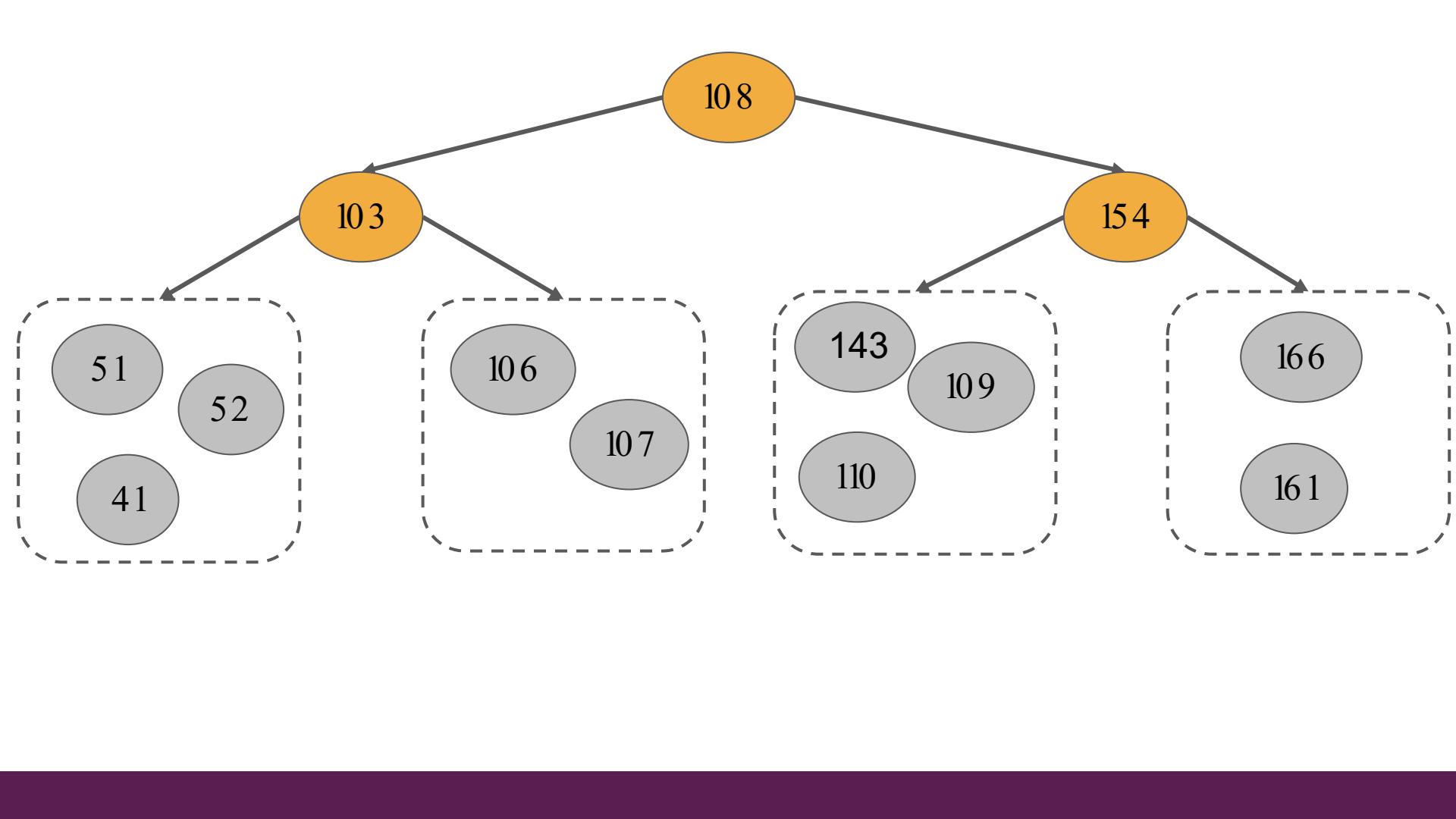


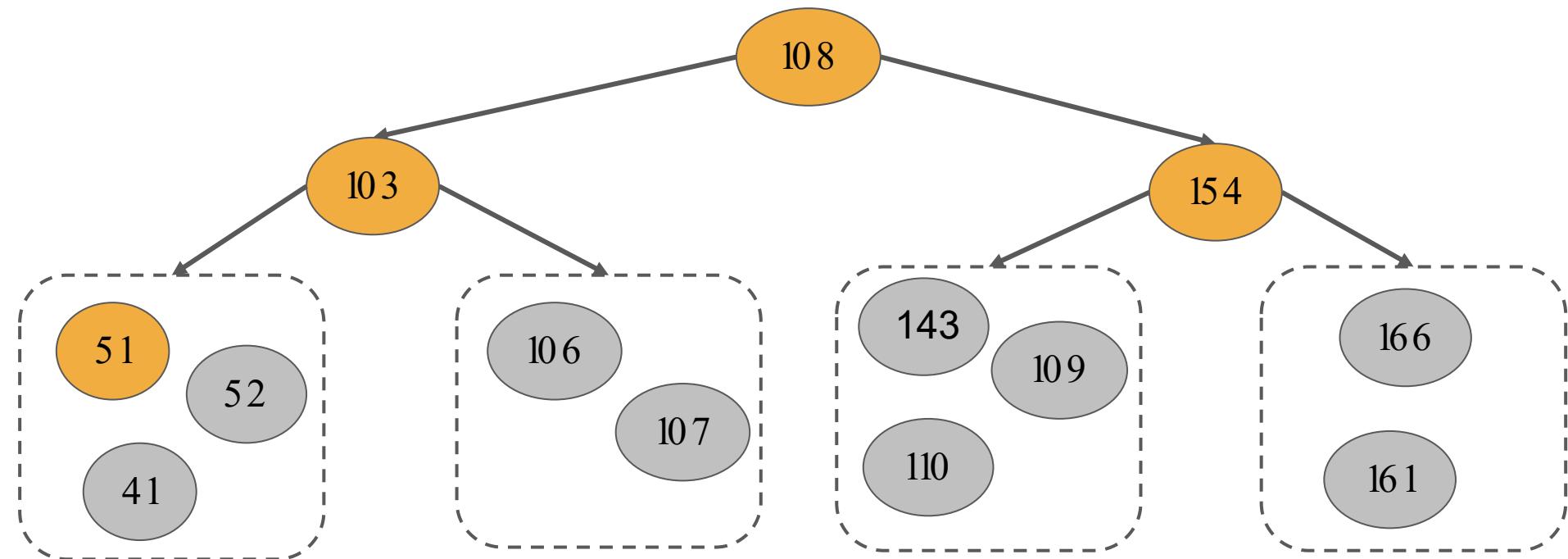


Pick the median element  
of the right side

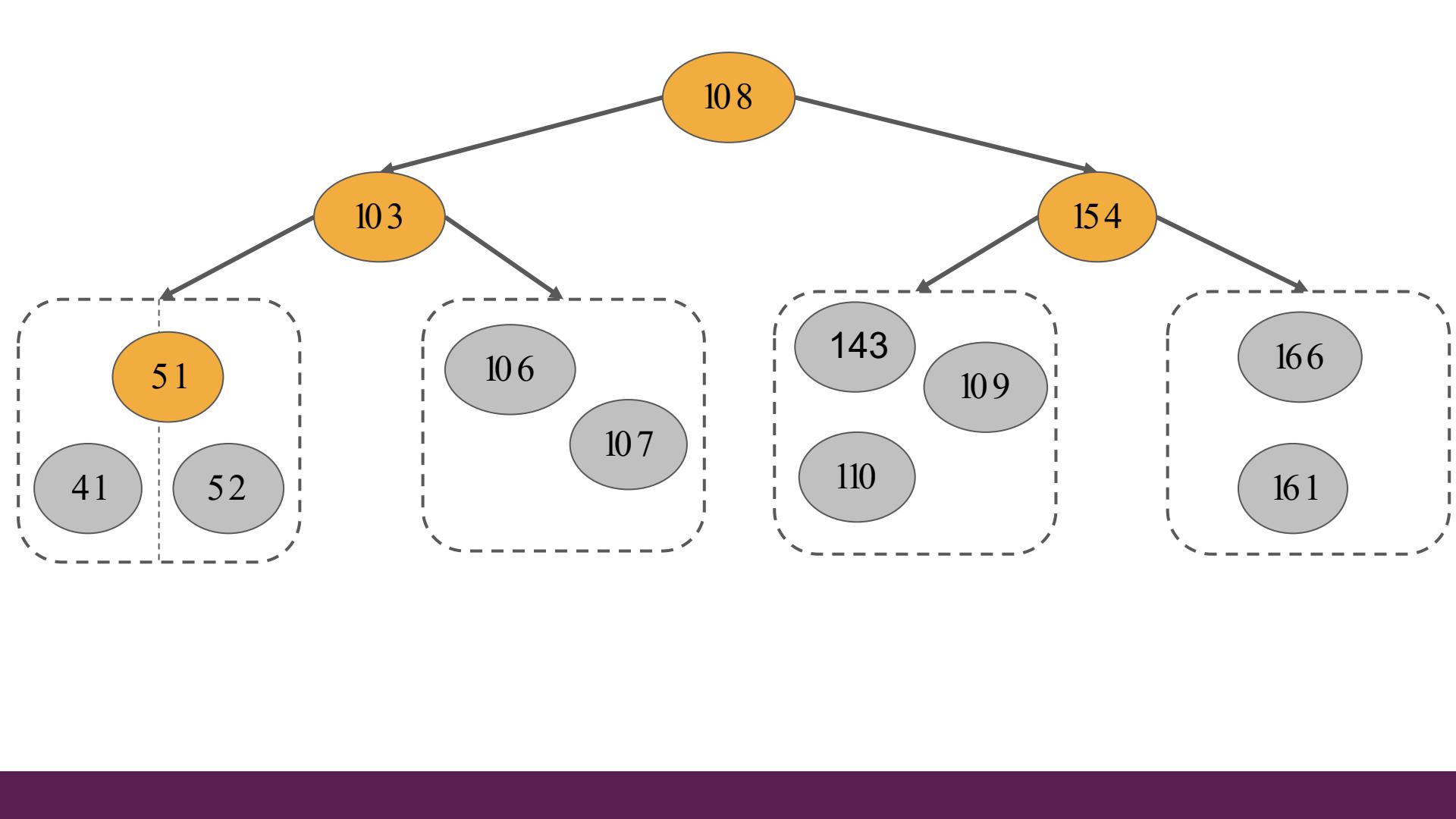


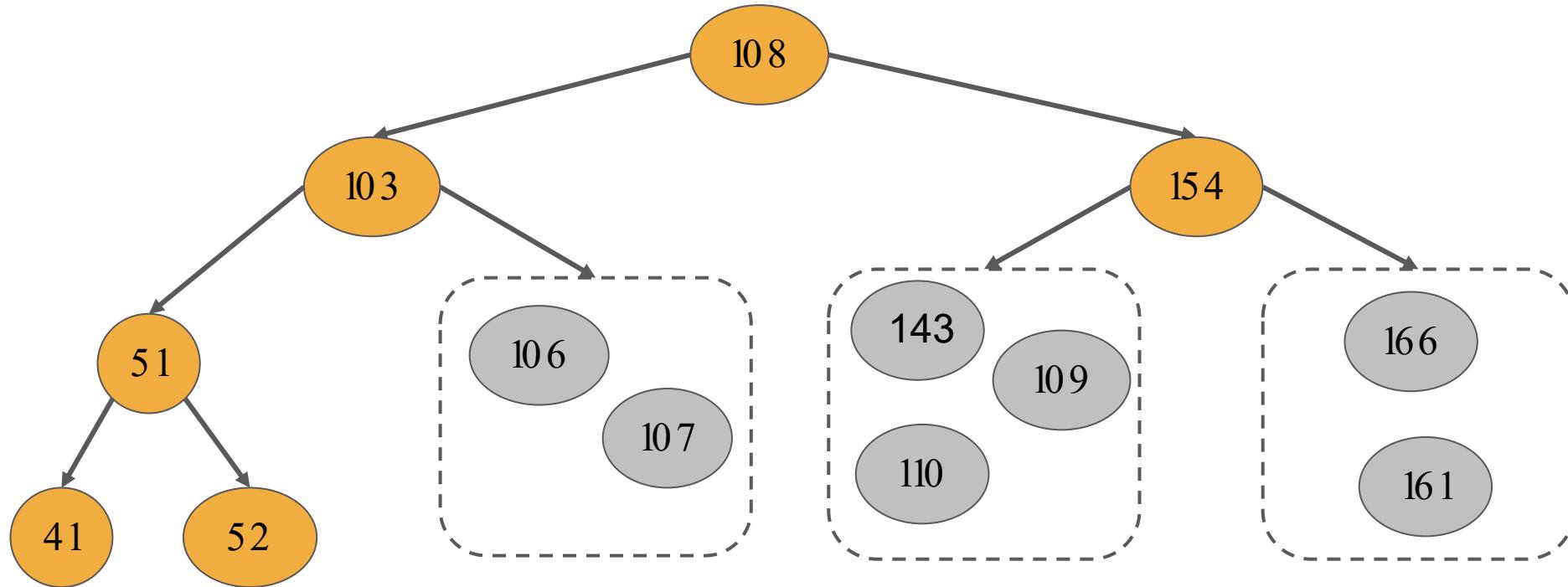


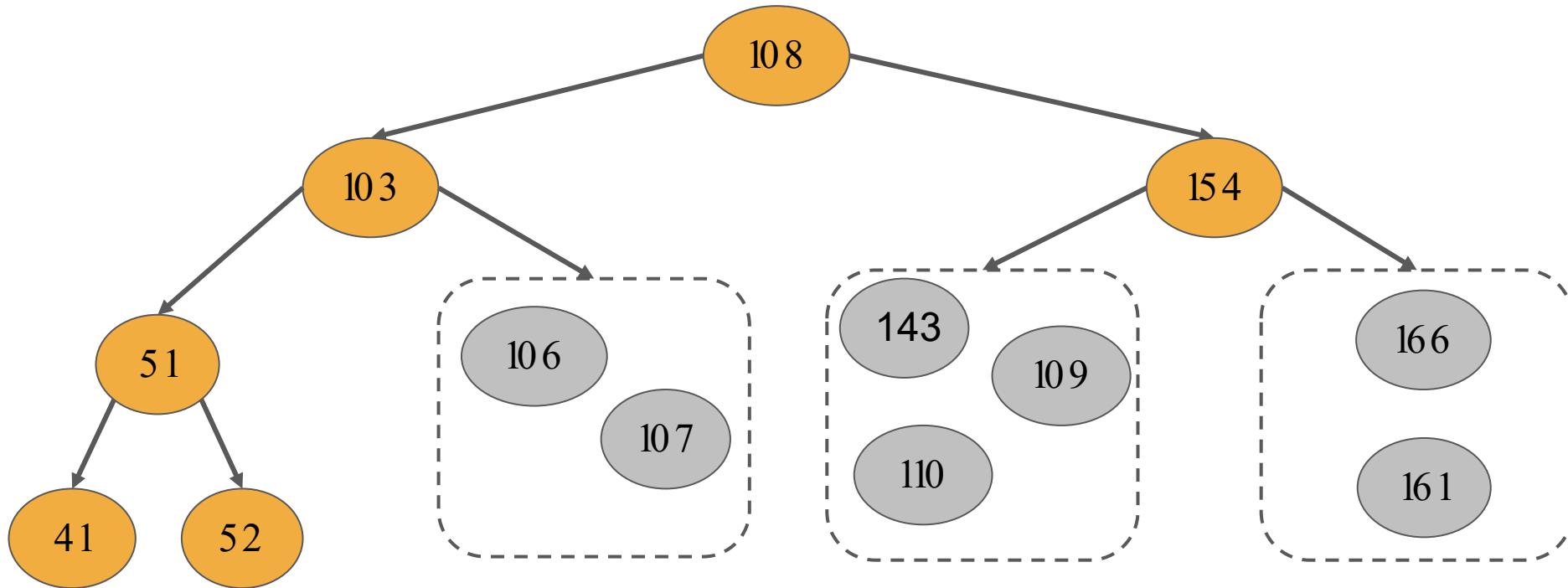




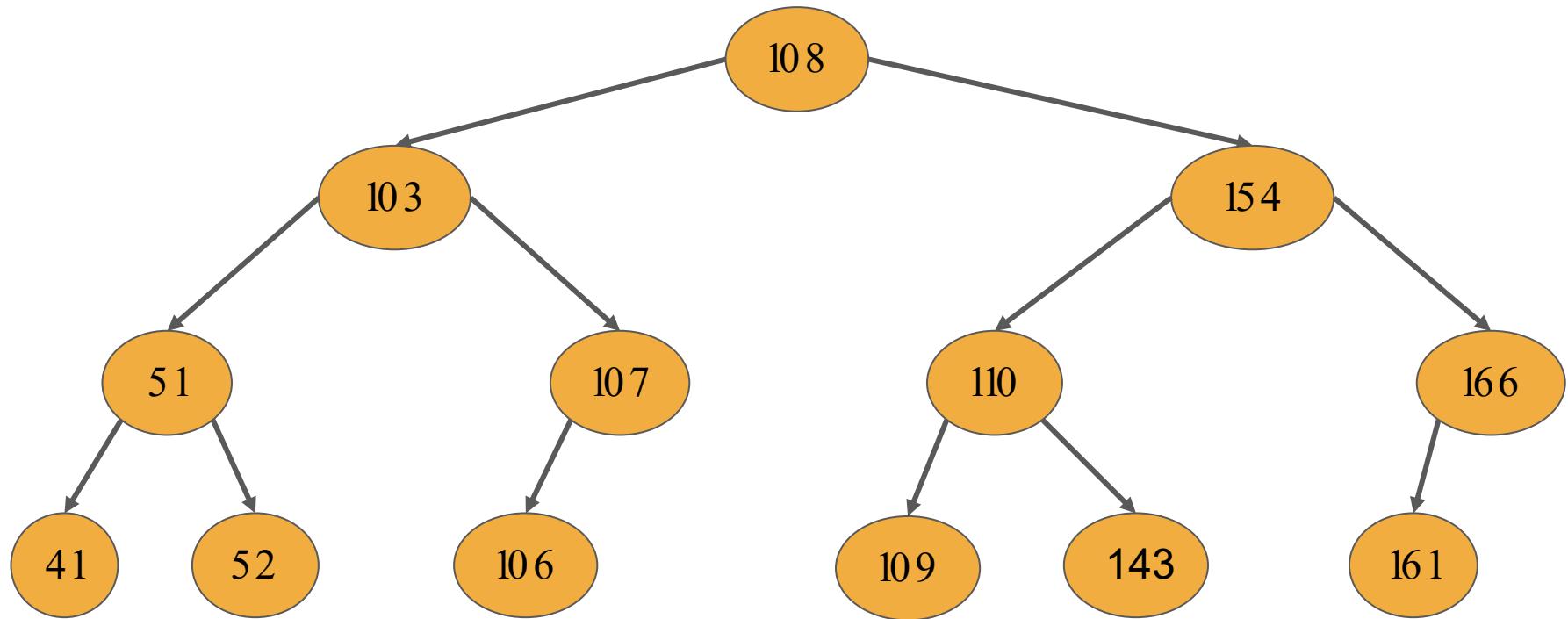
Pick the median element  
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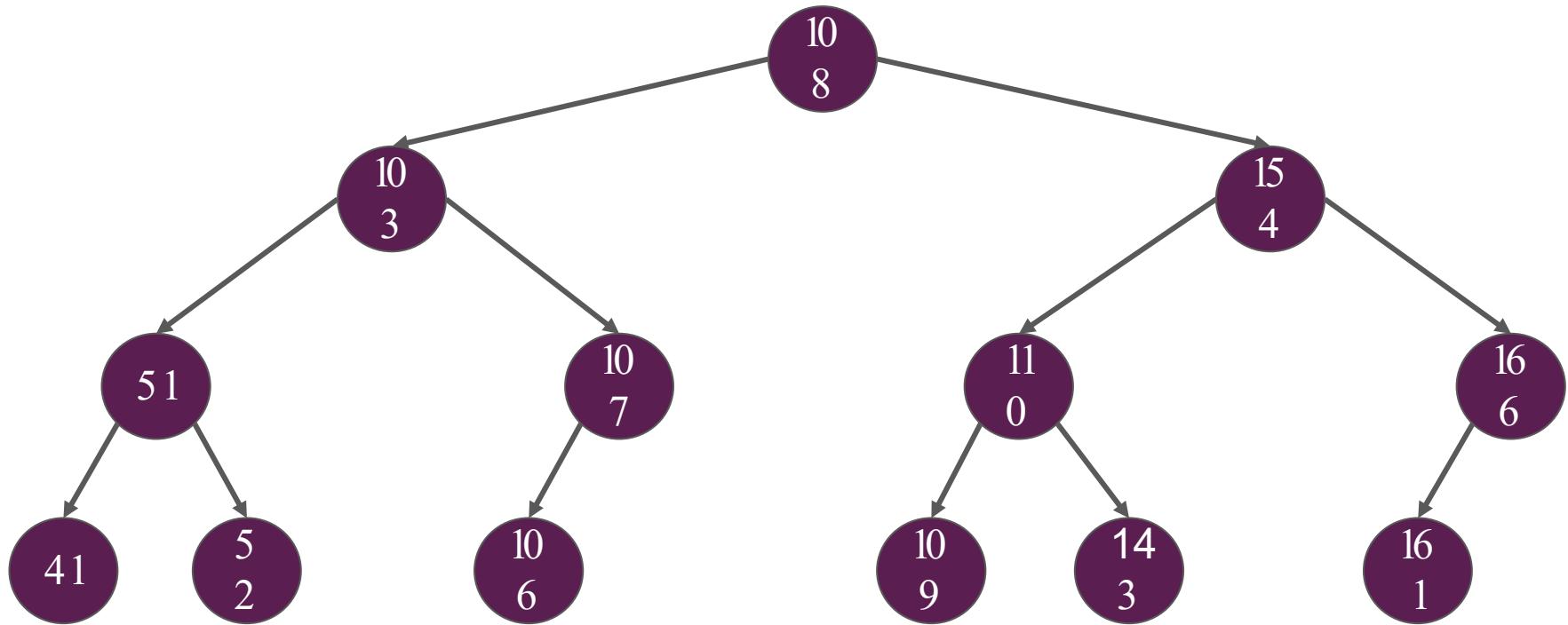




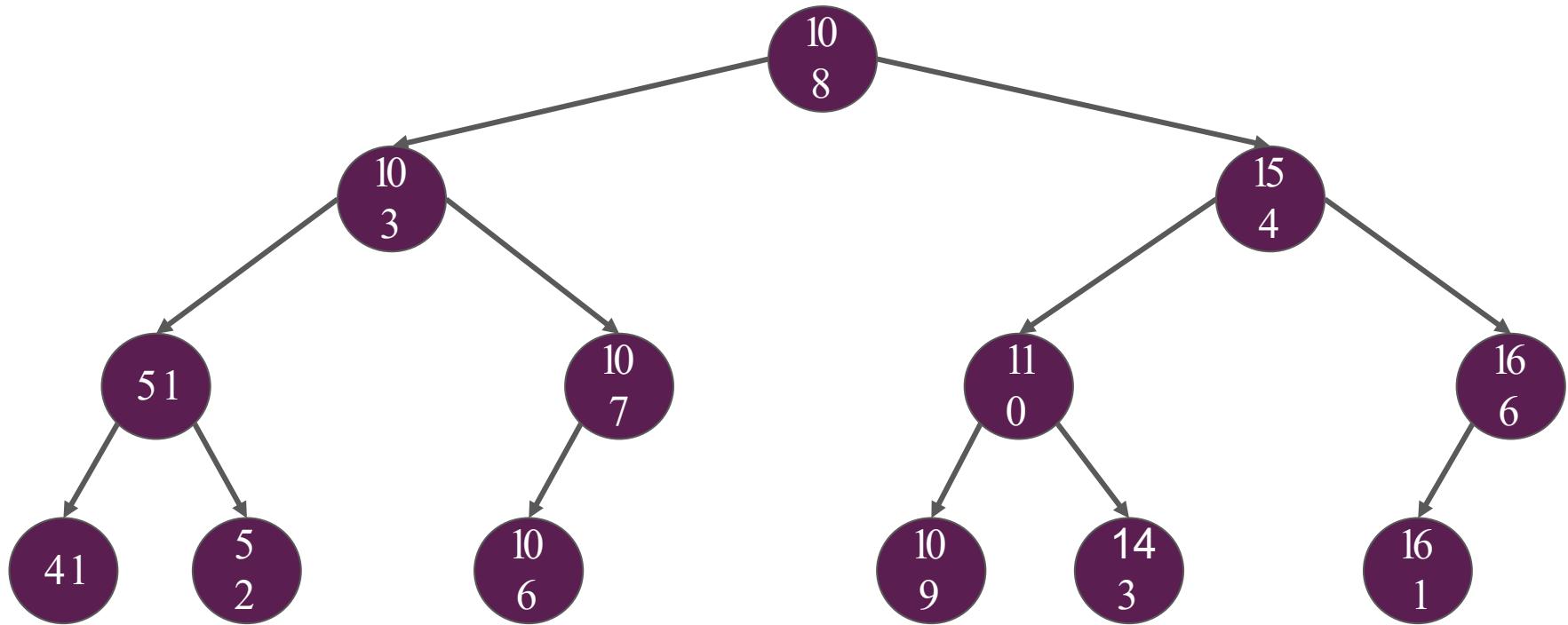


Keep repeating this process  
for all the subtrees!



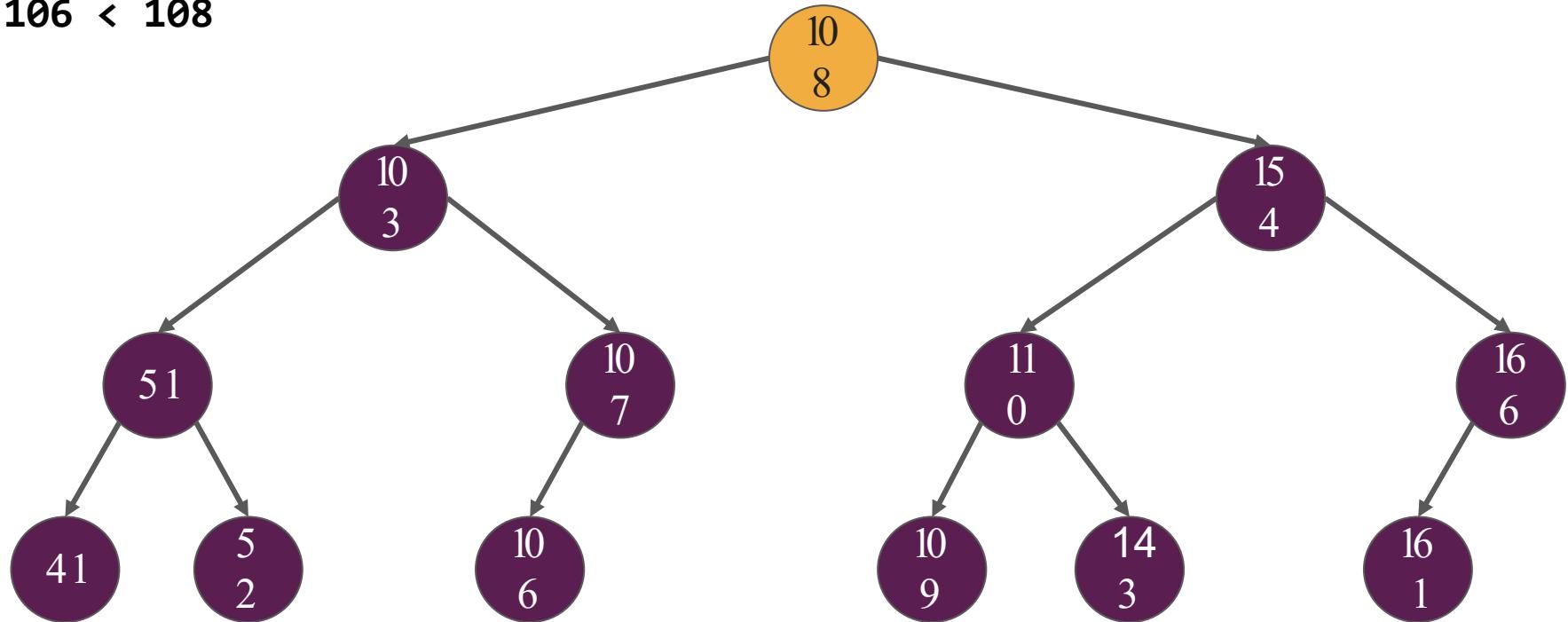


There are 13 nodes in the tree, but  
the path to each node is short  
 $(\sim O(\log 13))!$



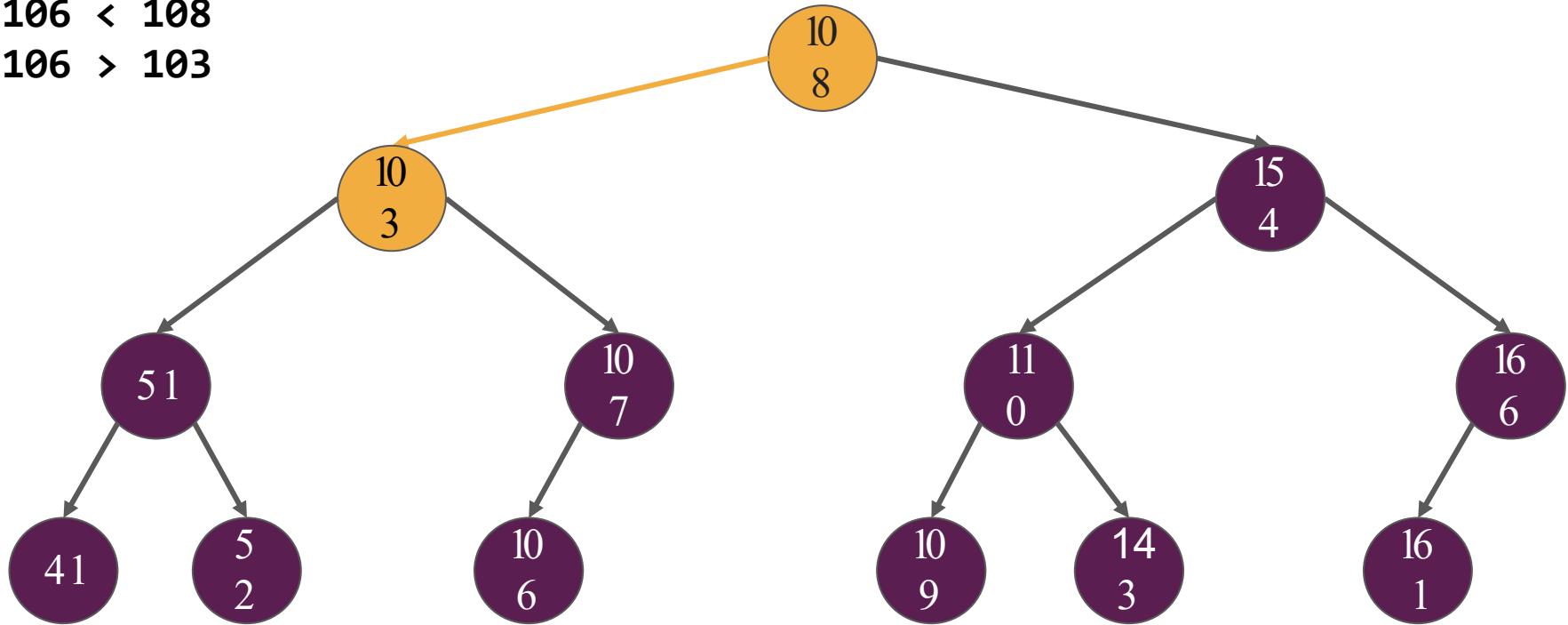
How could we check if **106** is in this tree?

**106 < 108**



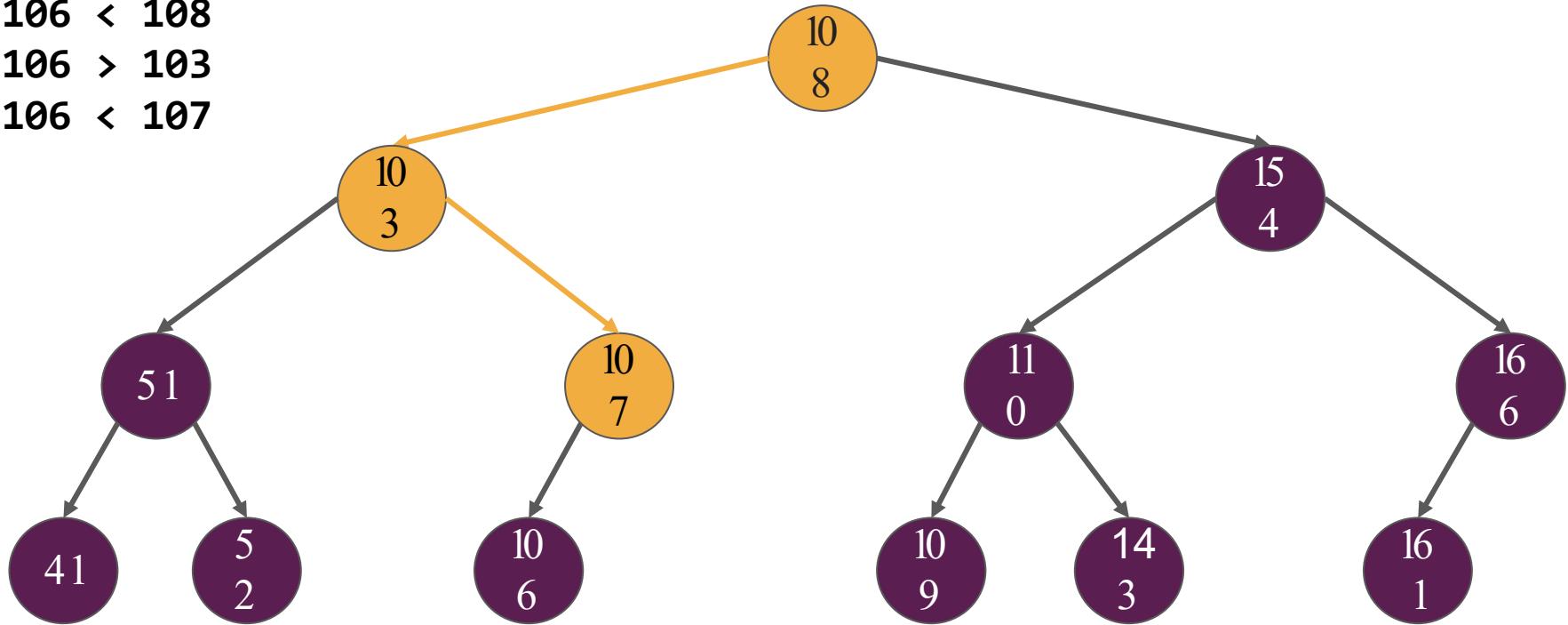
How could we check if **106** is in this tree?

**106 < 108**  
**106 > 103**



How could we check if **106** is in this tree?

106 < 108  
106 > 103  
106 < 107

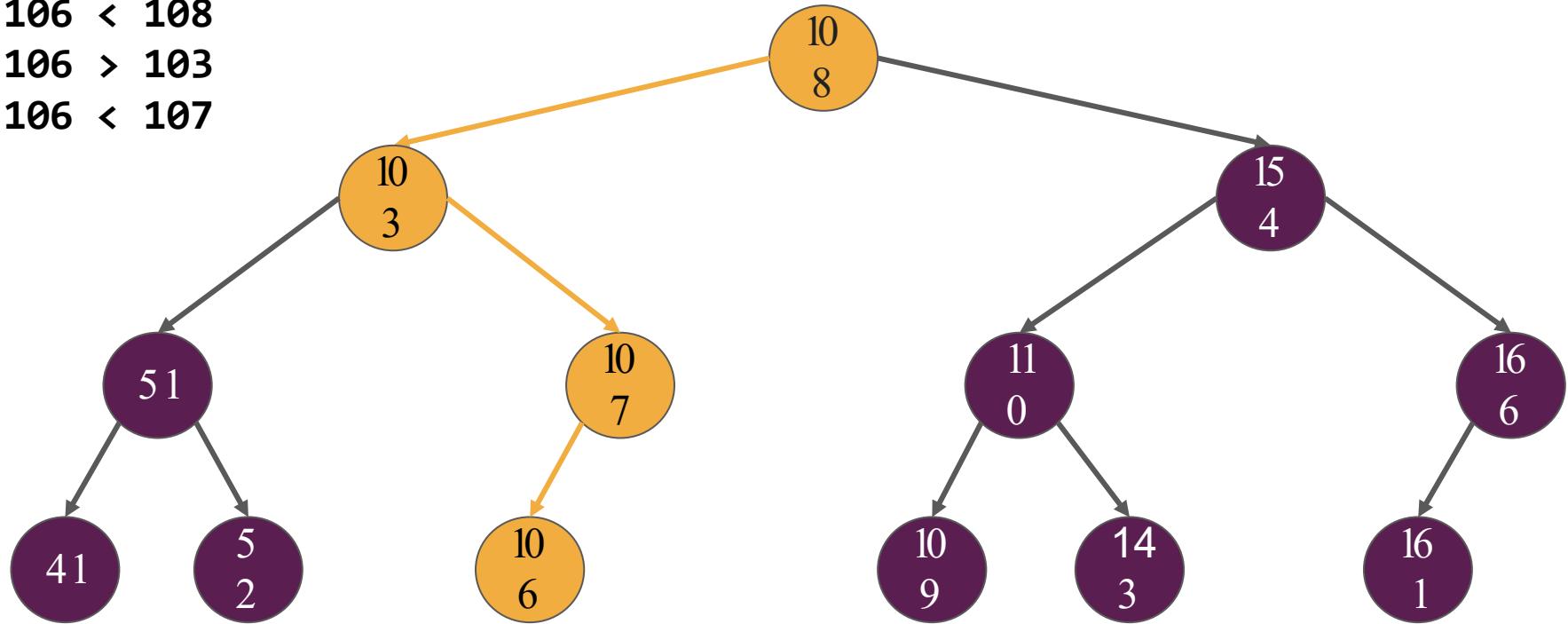


How could we check if **106** is in this tree?

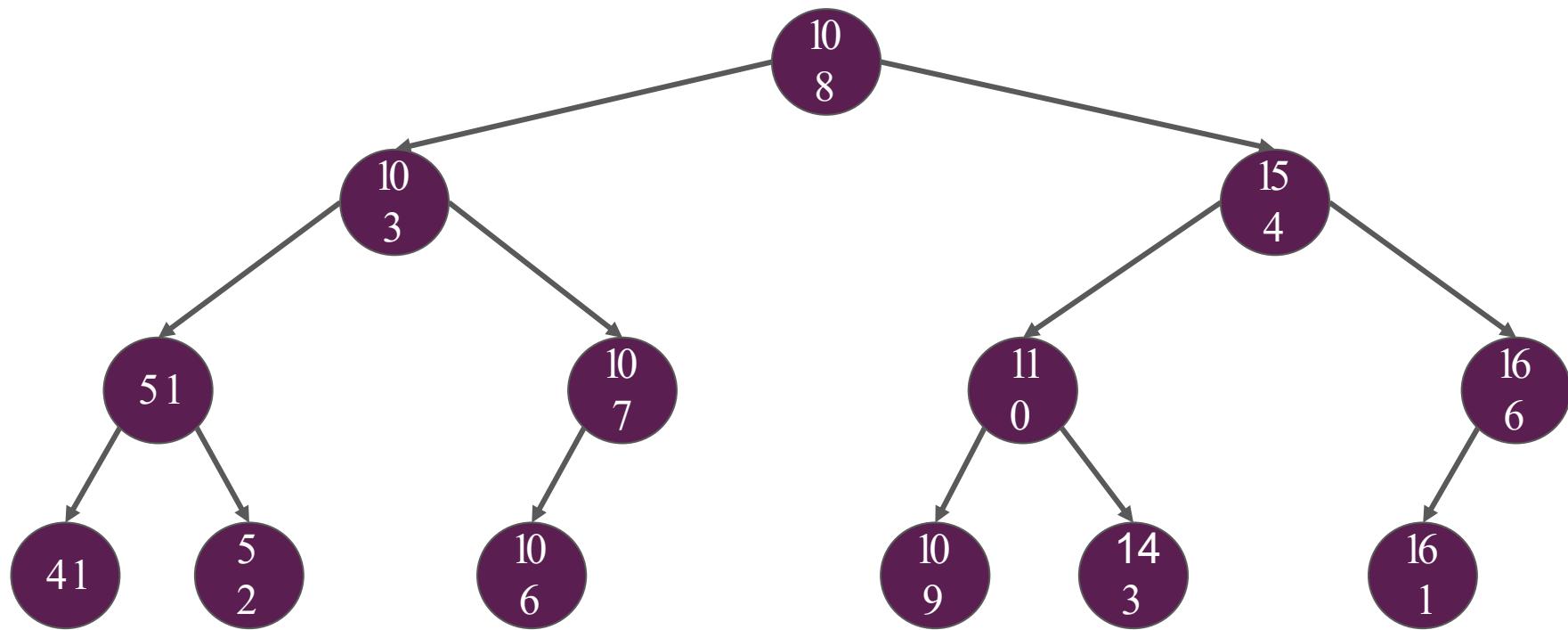
**106 < 108**

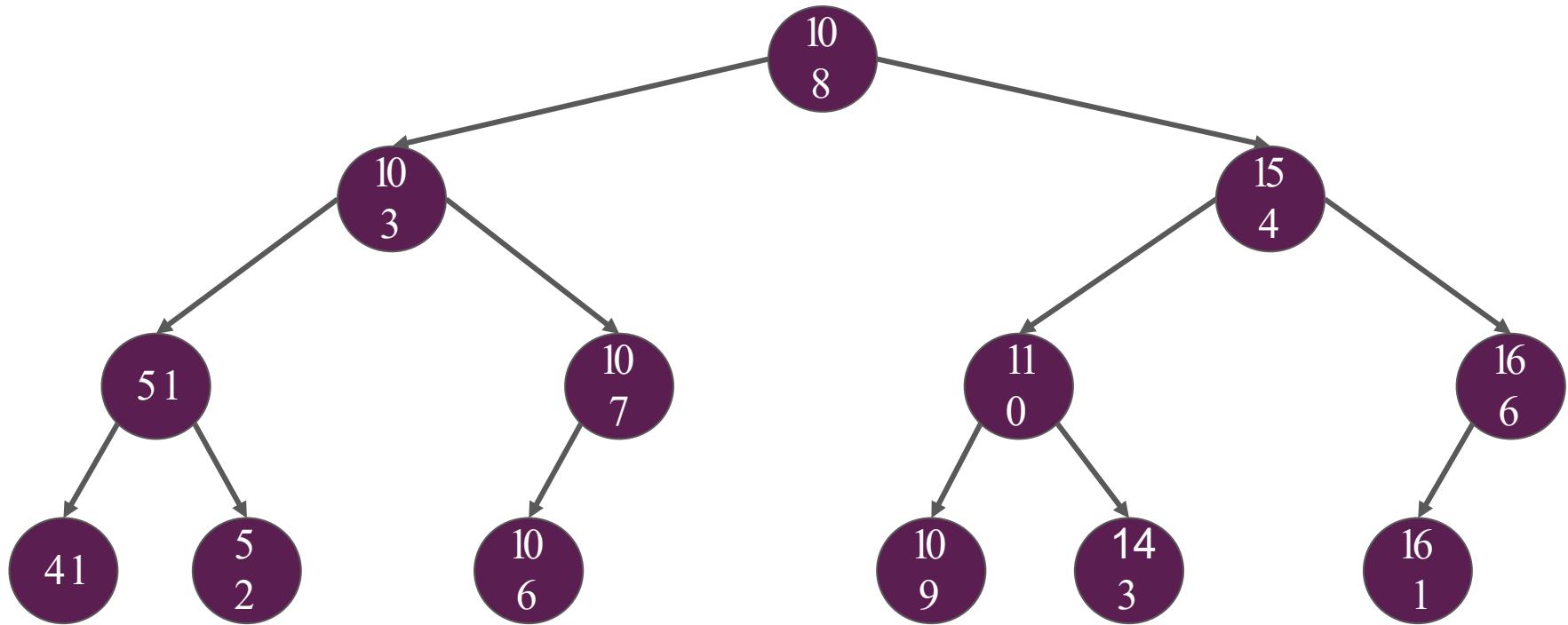
**106 > 103**

**106 < 107**



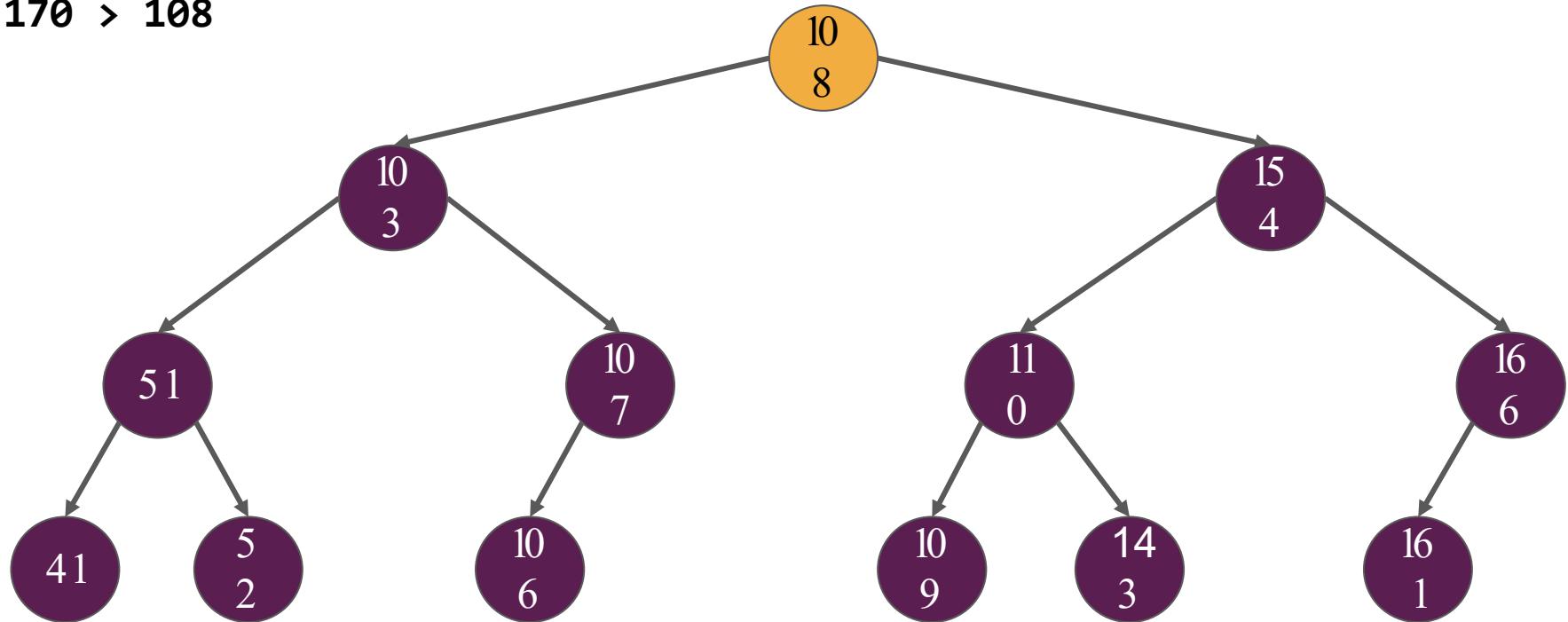
We found **106** so we're done!





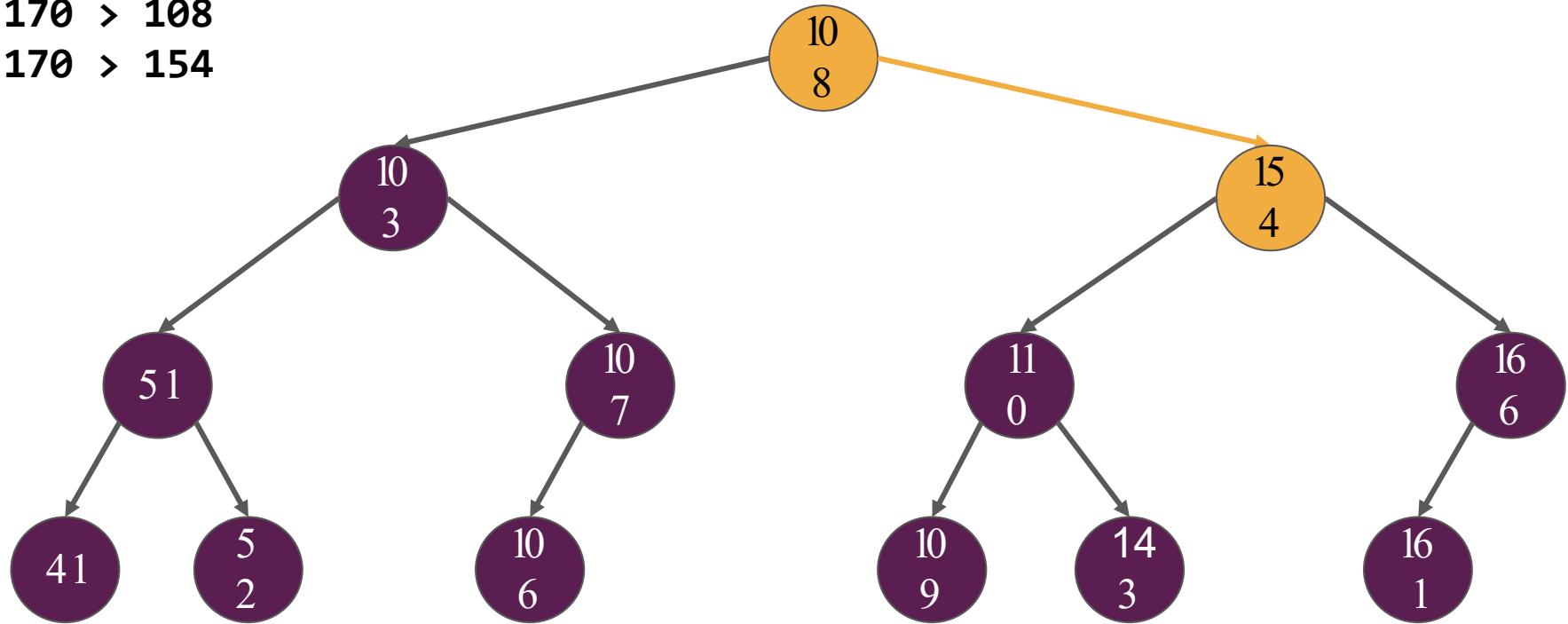
How could we check if **170** is in this tree?

**170 > 108**



How could we check if **170** is in this tree?

170 > 108  
170 > 154

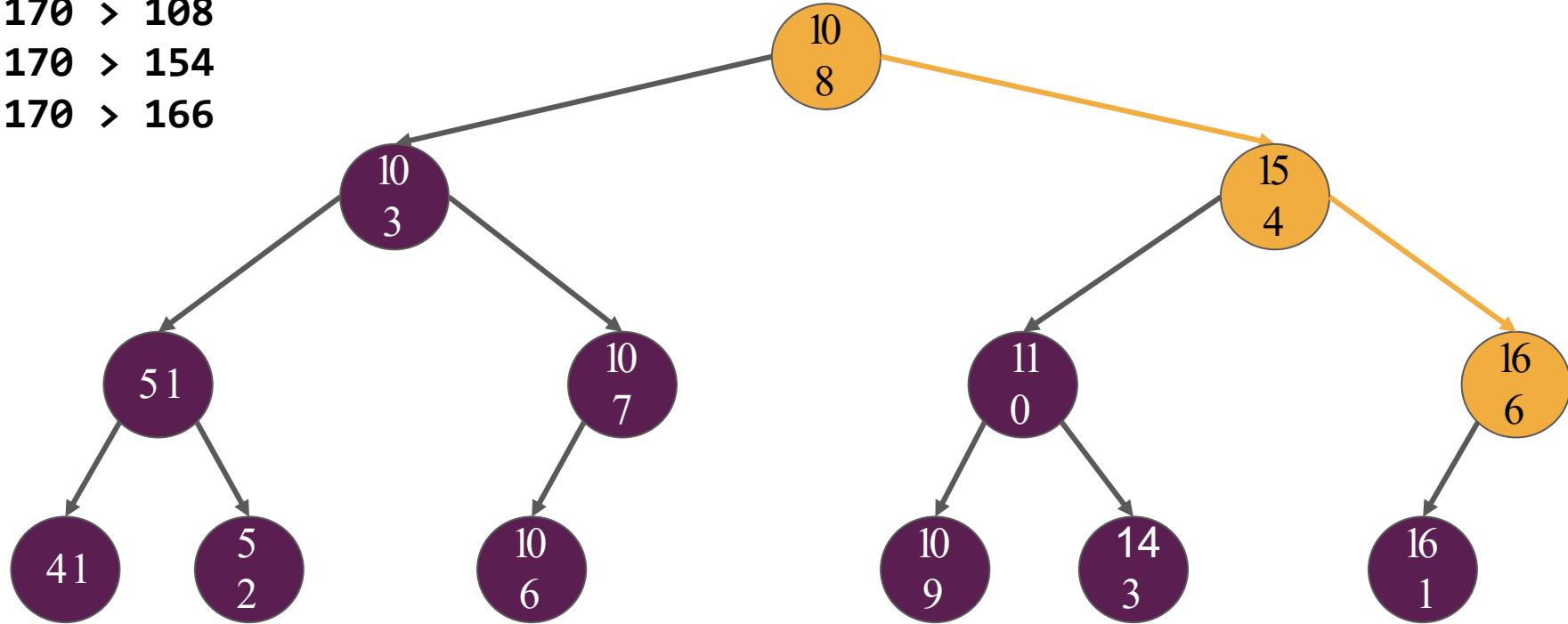


How could we check if **170** is in this tree?

170 > 108

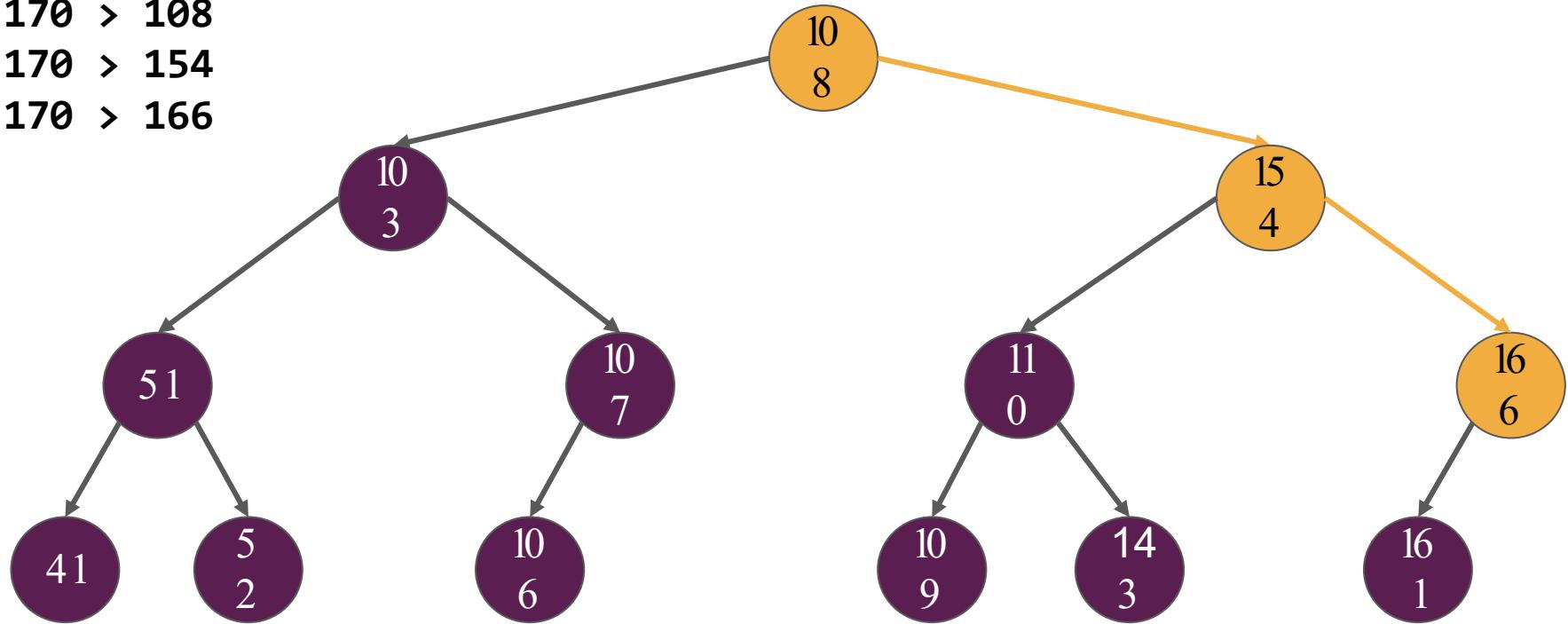
170 > 154

170 > 166



How could we check if **170** is in this tree?

170 > 108  
170 > 154  
170 > 166



Right child is **nullptr** so we're  
done!

# Building a BST

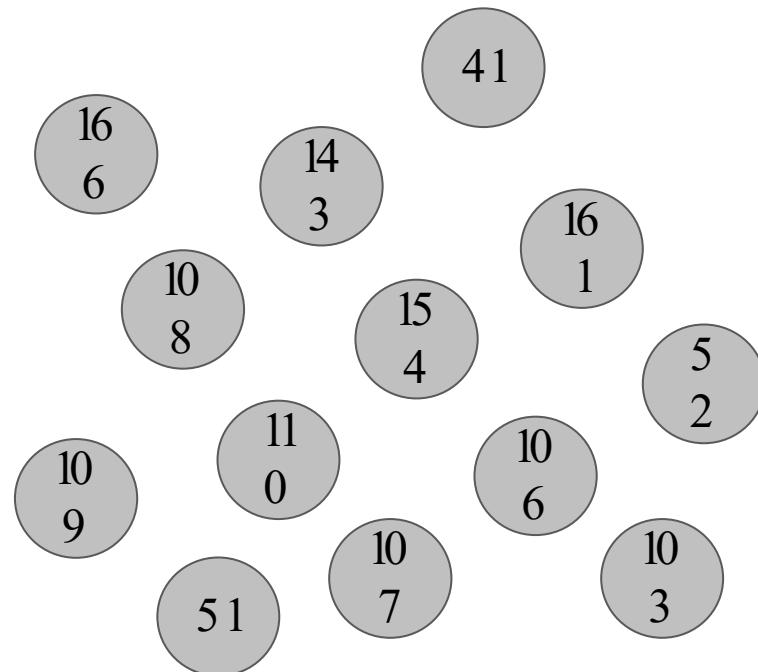
- An **optimal BST** is built by repeatedly choosing the median element as the root node of a given subtree and then separating elements into groups less than and greater than that median.

# Building a BST

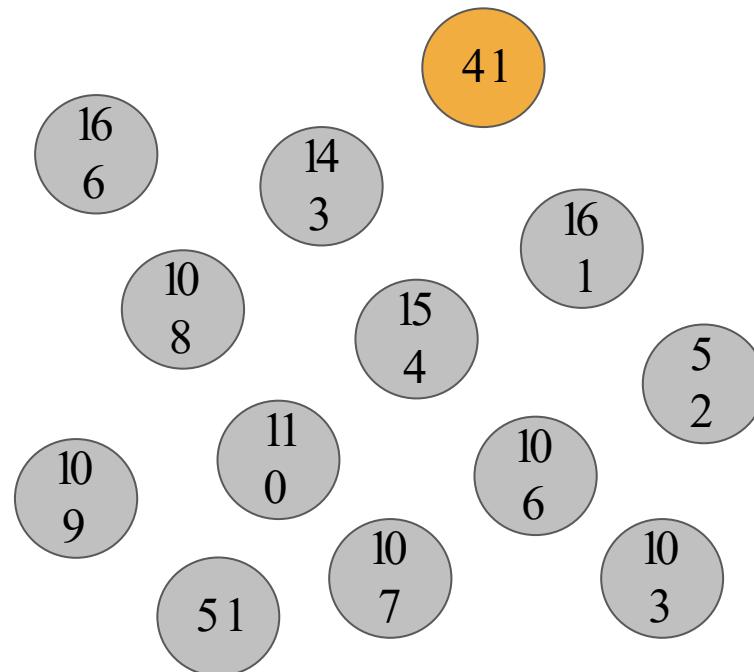
- An **optimal BST** is built by repeatedly choosing the median element as the root node of a given subtree and then separating elements into groups less than and greater than that median.

What does “optimal” mean?

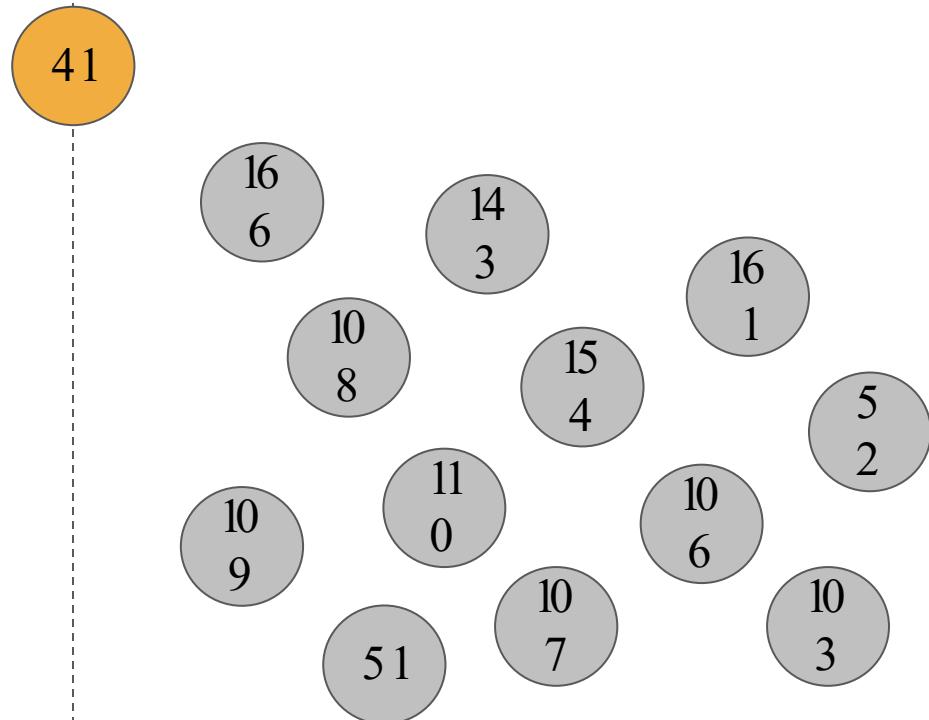
# What if we didn't choose the median?



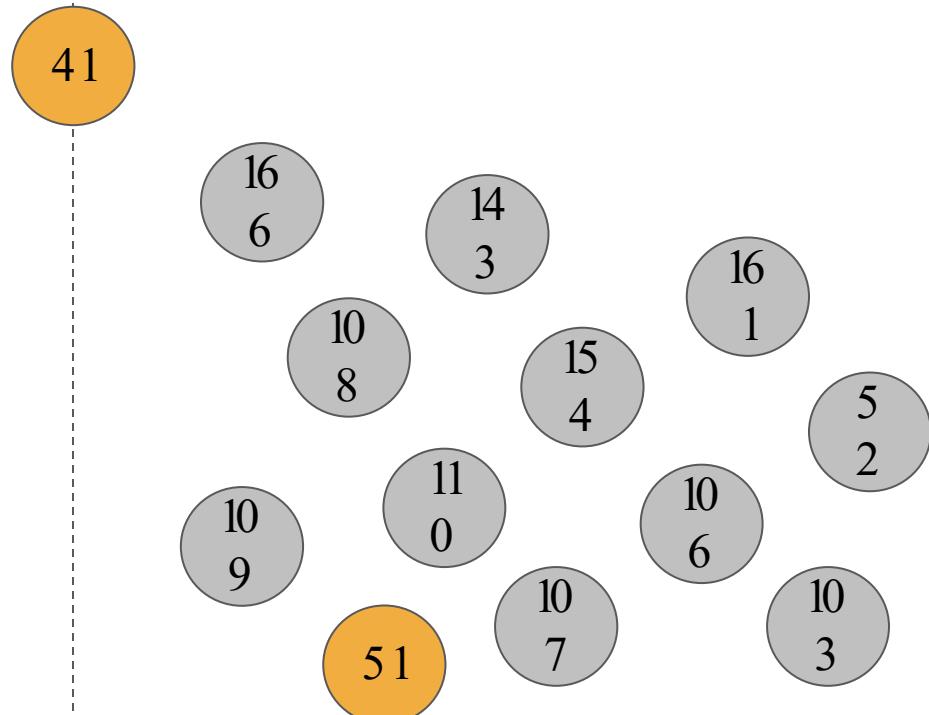
Let's choose the smallest element instead...



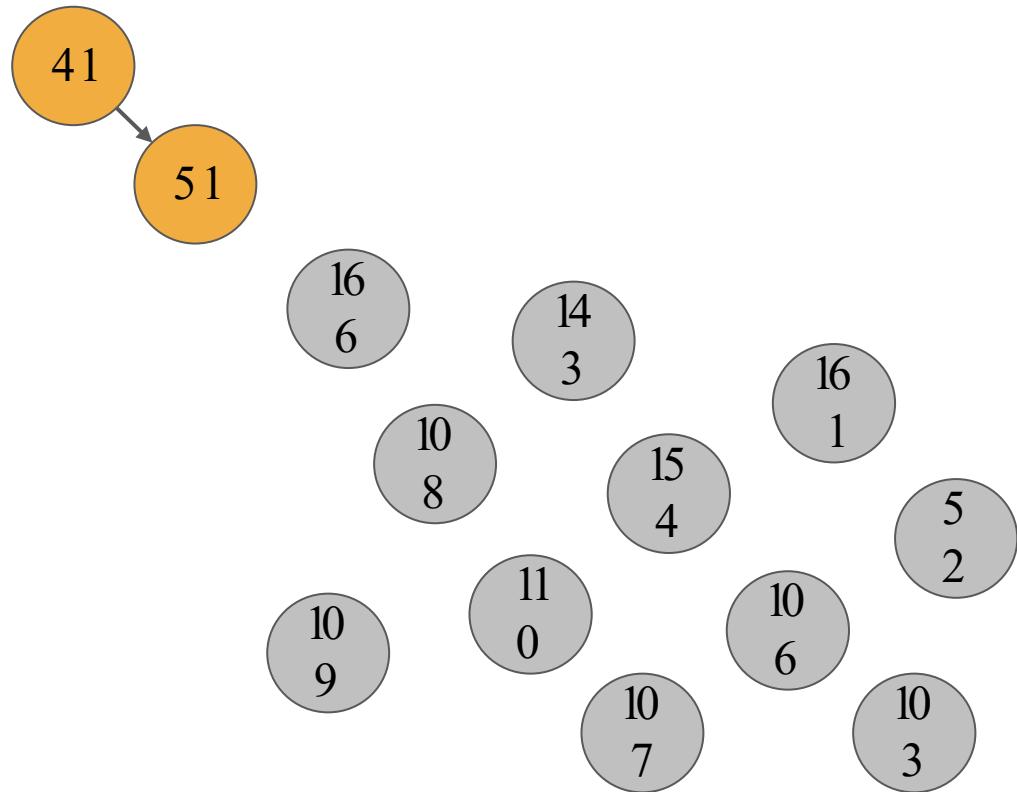
Let's choose the smallest element instead...

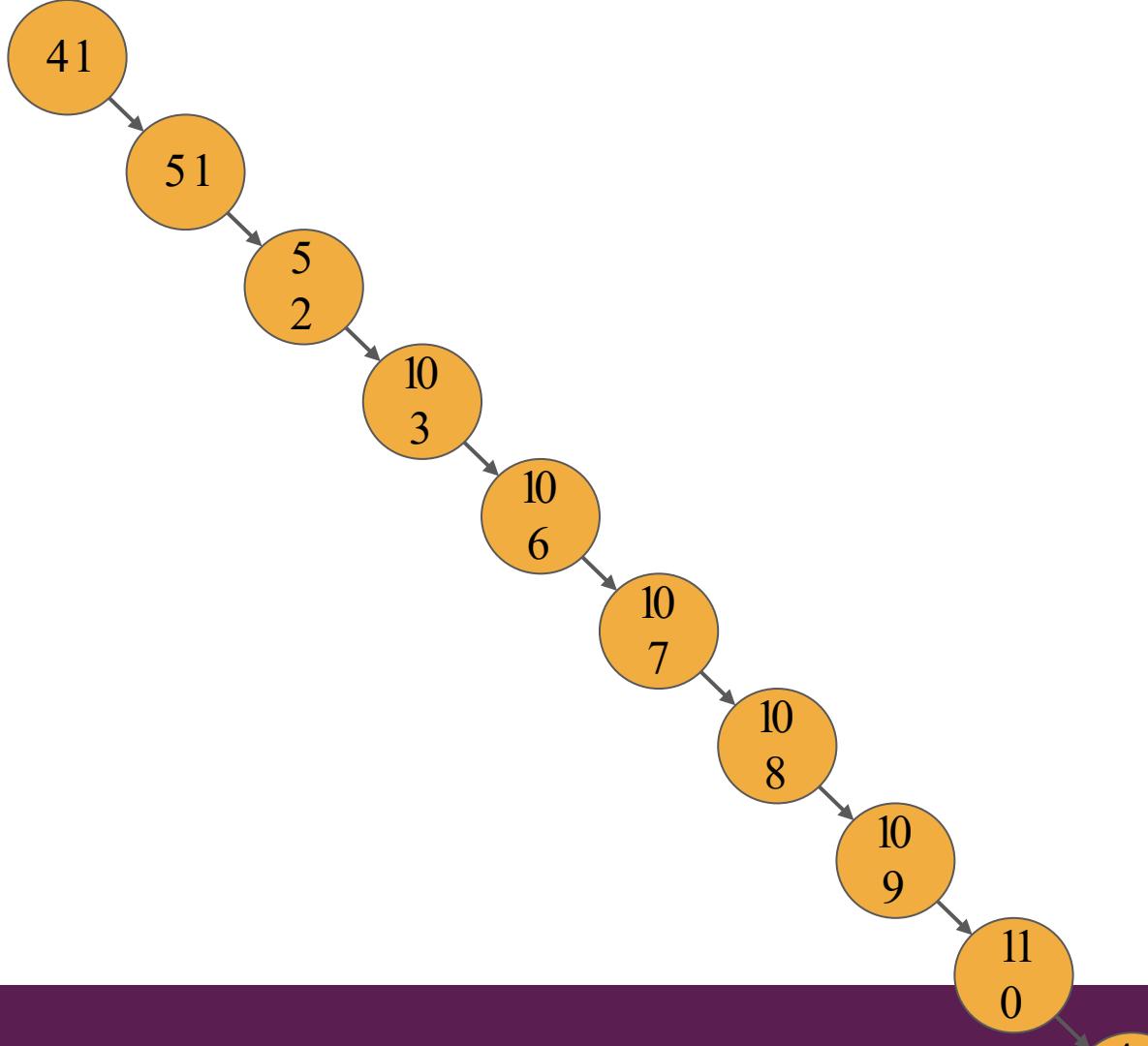


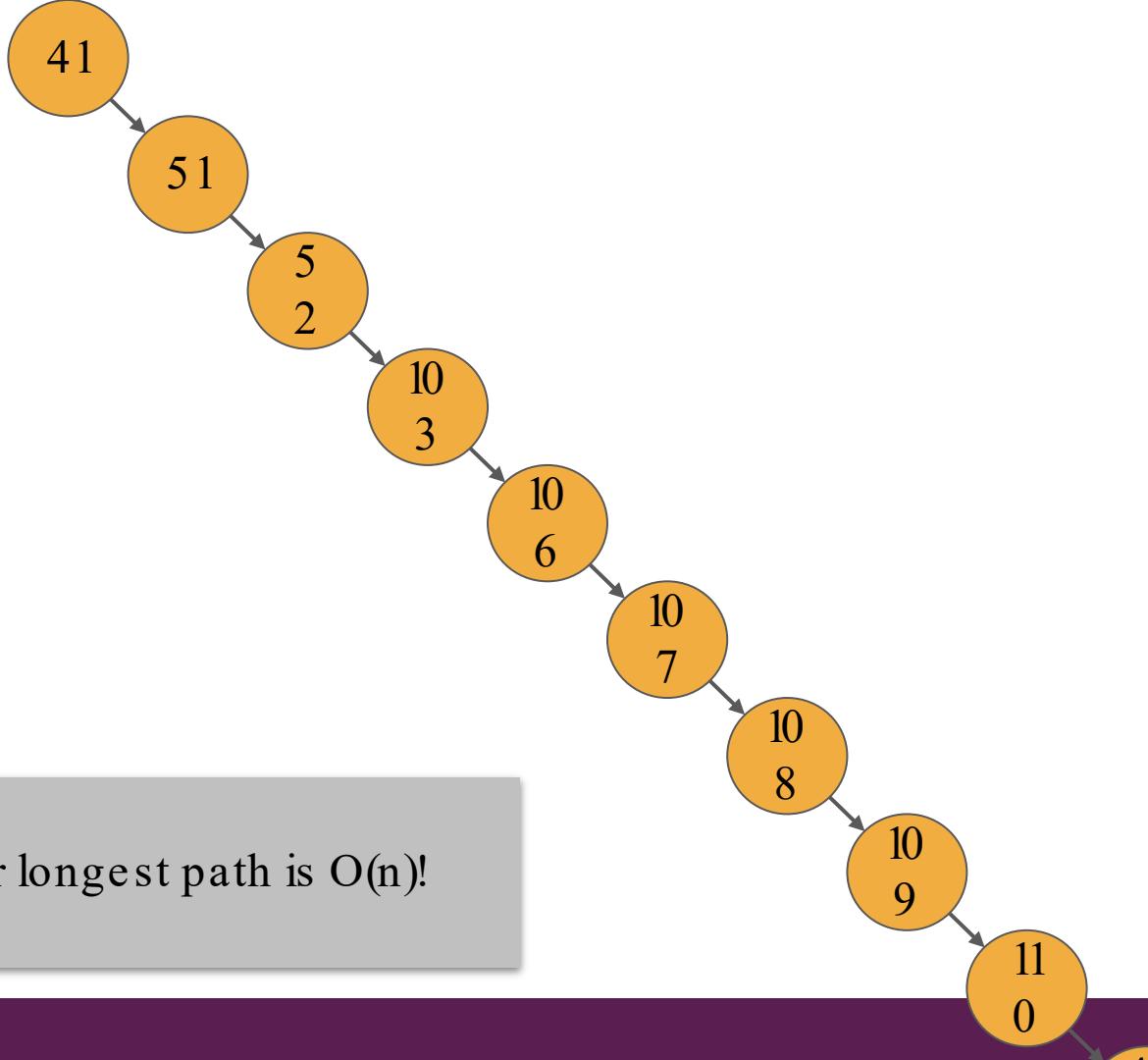
Let's choose the smallest element instead...



Let's choose the smallest element instead...



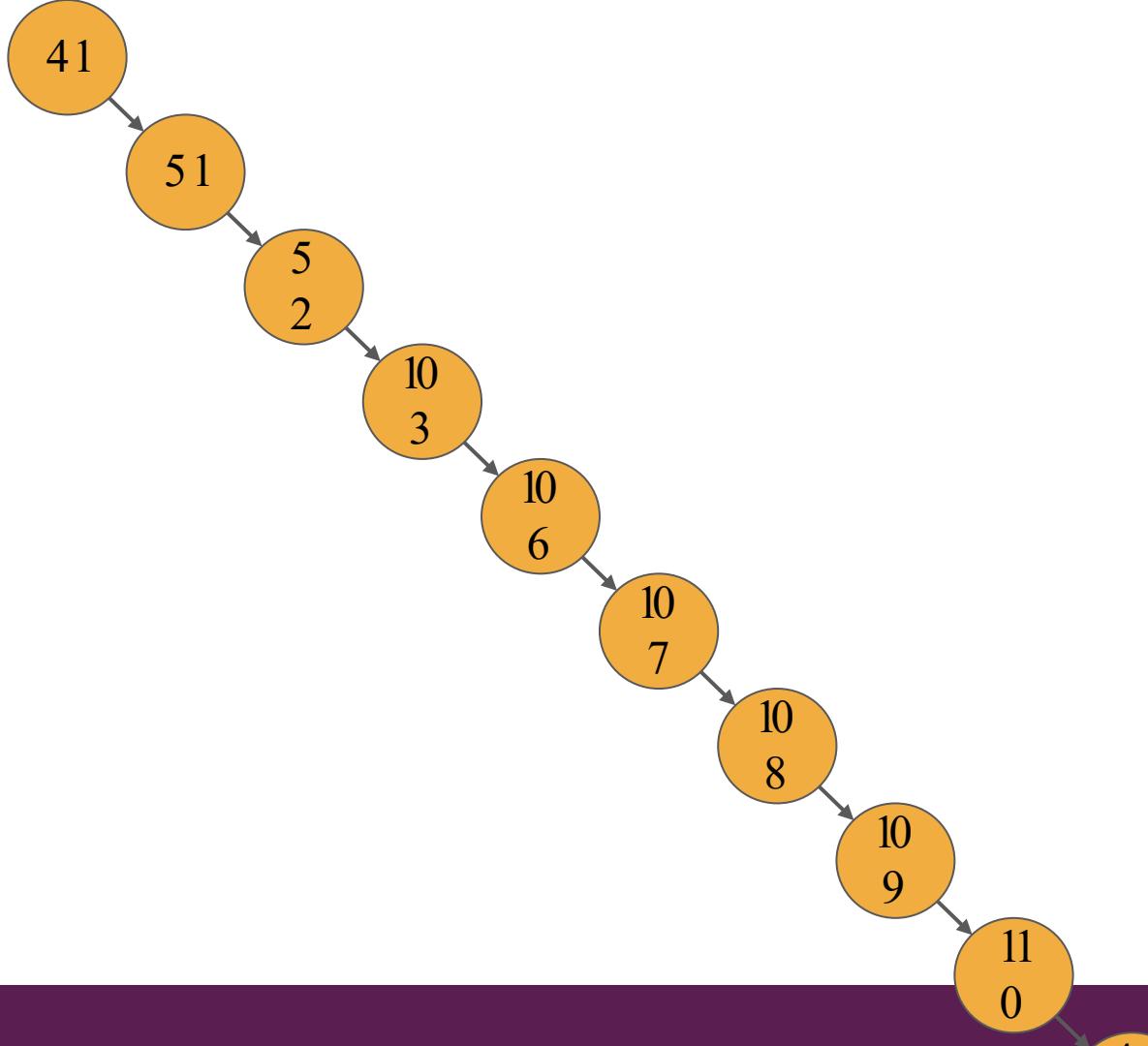




Now our longest path is  $O(n)$ !

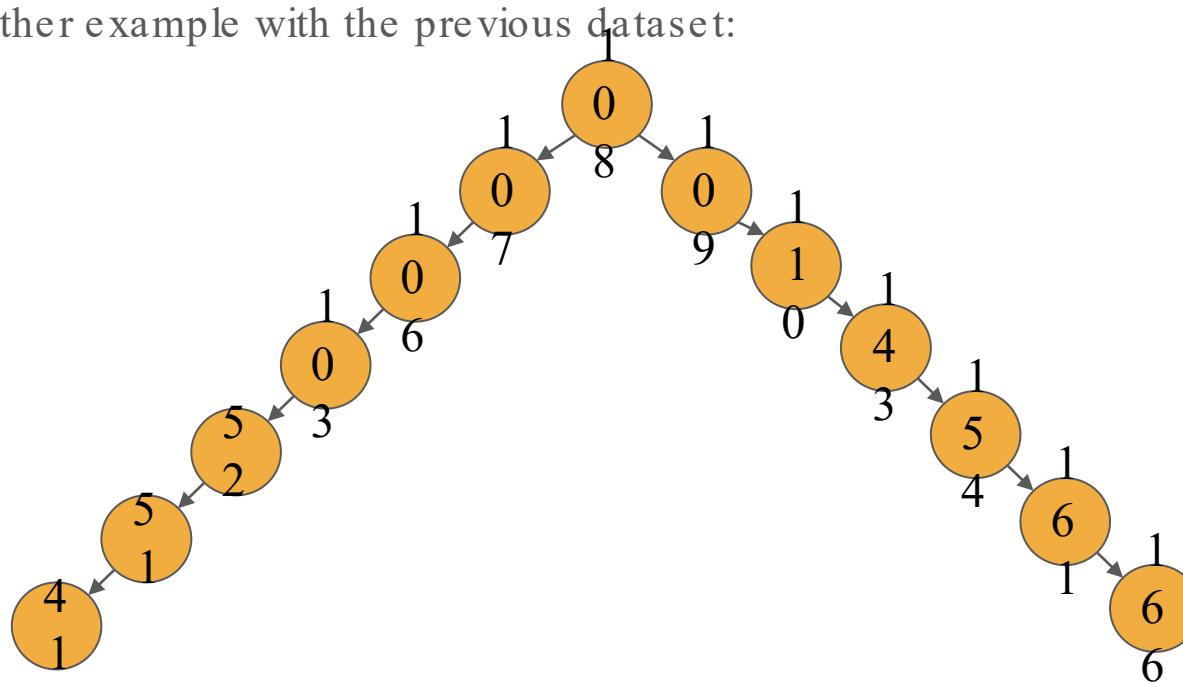
# Takeaways

- There are multiple valid BSTs for the same set of data.



# Takeaways

- There are multiple valid BSTs for the same set of data.
  - Another example with the previous dataset:



# Takeaways

- There are multiple valid BSTs for the same set of data.
- How you construct the tree/the order in which you add the elements to the tree matters!

# Takeaways

- There are multiple valid BSTs for the same set of data.
- How you construct the tree/the order in which you add the elements to the tree matters!
- A binary search tree is **balanced** if its height is  **$O(\log n)$** , where  **$n$**  is the number of nodes in the tree (i.e. left/right subtrees don't differ in height by more than 1).
  - Lookup, insertion, and deletion with balanced BSTs all operate in  **$O(\log n)$**  runtime.

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  - **Theorem:** If you start with an empty tree and add in random values, then with high probability the tree is balanced. → take CS161 to learn why!
  - A **self-balancing** BST reshapes itself on insertions and deletions to stay balanced (how to do this is beyond the scope of this class).

# Announcements

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- Assignment 5 is due **tonight at 11:59pm PDT**.
- Assignment 6 will be released by the end of the day tomorrow and will be due on **Wednesday, August 11 at 11:59pm PDT** This is a hard deadline – there is **no grace period and no submissions will be accepted after this time** .
- The End-quarter Assessment will take place over 3 days from **Friday, August 13 to Sunday, August 15**.

# Implementing Sets with BSTs

# We're going to implement a Set using a BST!

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```
struct TreeNode {  
    std::string data;  
    TreeNode* left;  
    TreeNode* right;  
  
    // default constructor does not initialize  
    TreeNode() {}  
    // 3-arg constructor sets fields from arguments  
    TreeNode(std::string d, TreeNode* l, TreeNode* r) {  
        data = d;  
        left = l;  
        right = r;  
    }  
};
```

# We're going to implement a Set using a BST!

- Our Set will only store strings as its data type
- We have a header file that will include a public interface already defined.

# OurSet Public Interface

```
class OurSet {  
public:  
    OurSet();    // constructor  
    ~OurSet();   // destructor  
  
    bool contains(string value);  
    void add(string value);  
    void remove(string value);  
    void clear();  
    int size();  
    bool isEmpty();  
    void printSetContents();  
  
private:  
    /* To be defined soon! */  
};
```

# We're going to implement a Set using a BST!

- Our Set will only store strings as its data type
- We have a header file that will include a public interface already defined.
- As we write the Set methods, think about how their runtimes would change for a balanced vs. an unbalanced BST.
  - Note: Actual sets are self-balancing, but we won't go into the details of how to implement that!

# How do we design OurSet?

We must answer the following three questions:

1. Member functions: *What public interface should OurSet support? What functions might a client want to call?*
2. Member variables: *What private information will we need to store in order to keep track of the data stored in OurSet?*
3. Constructor: *How are the member variables initialized when a new instance of OurSet is created?*

# OurSet Public Interface

```
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};
```

# Let's code it!

constructor and destructor

# Let's code it!

`size()`, `isEmpty()`, `clear()`

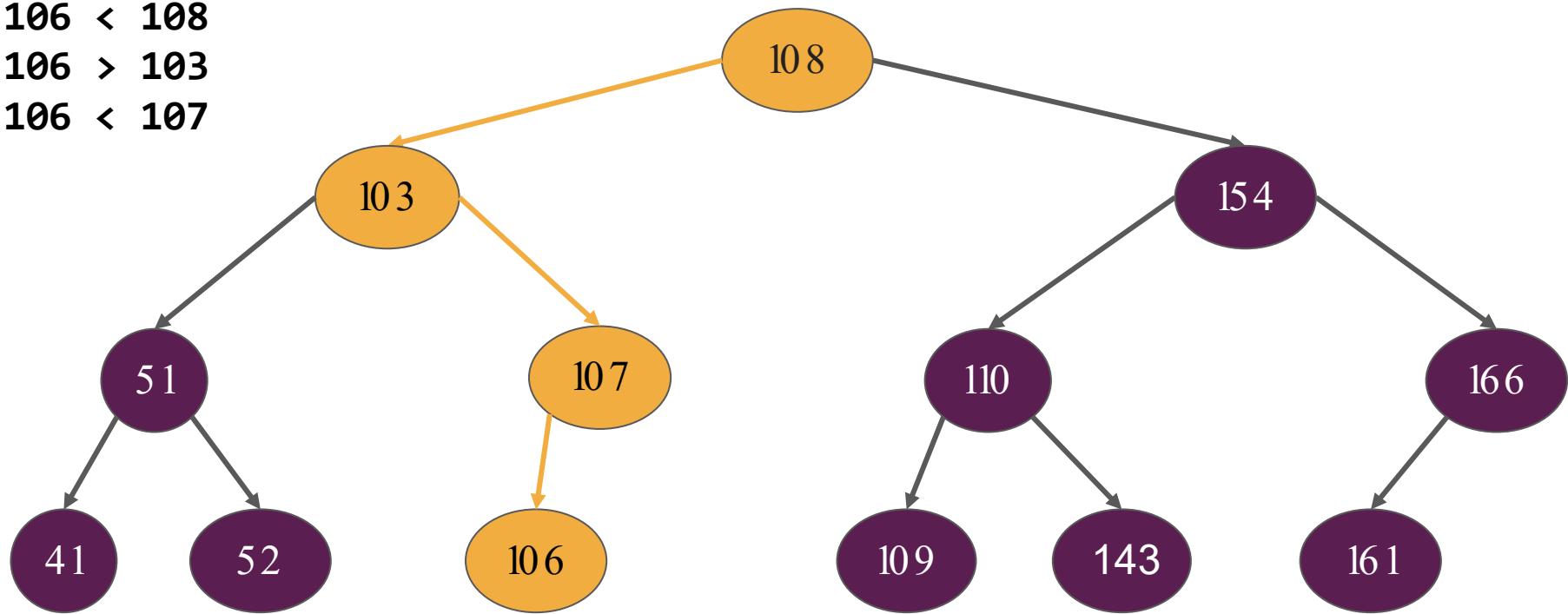
# Let's code it!

`printSetContents ()`

# OurSet Public Interface

```
class OurSet {  
public:  
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    void printSetContents();  
  
private:  
    /* To be defined soon! */  
};
```

106 < 108  
106 > 103  
106 < 107

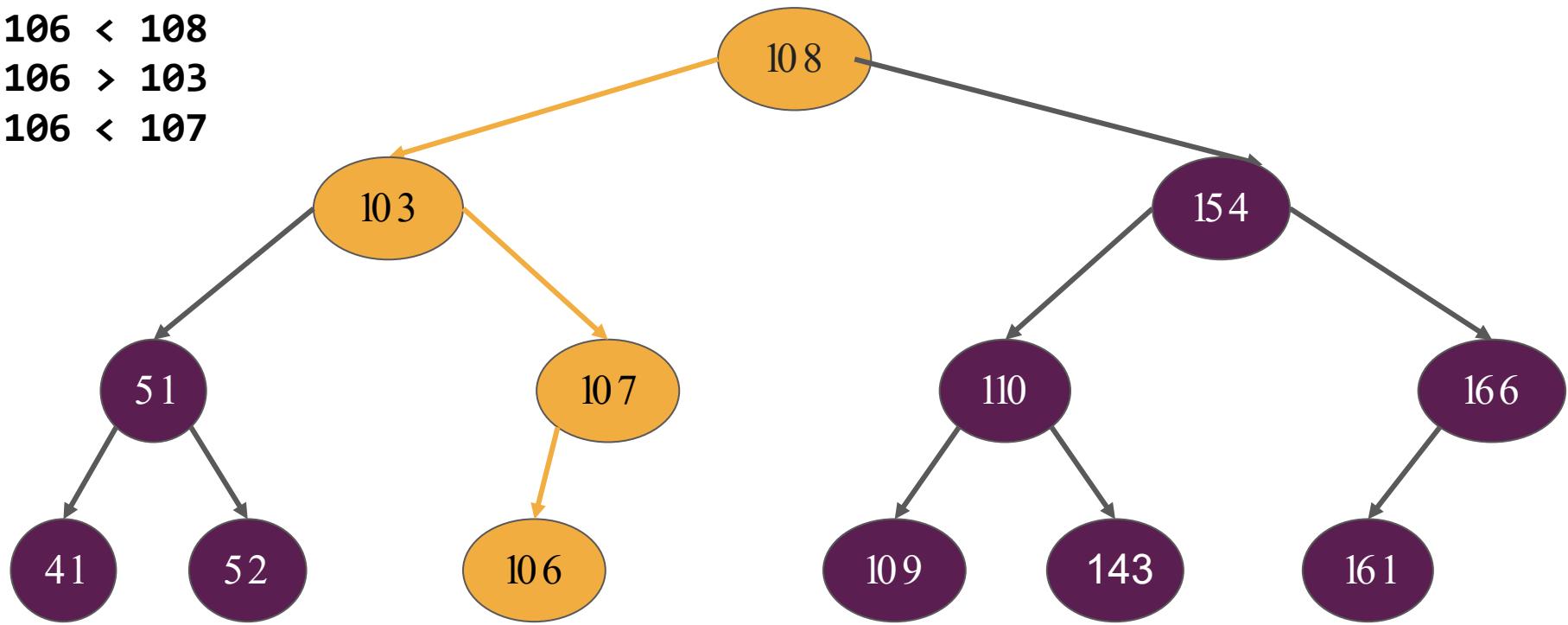


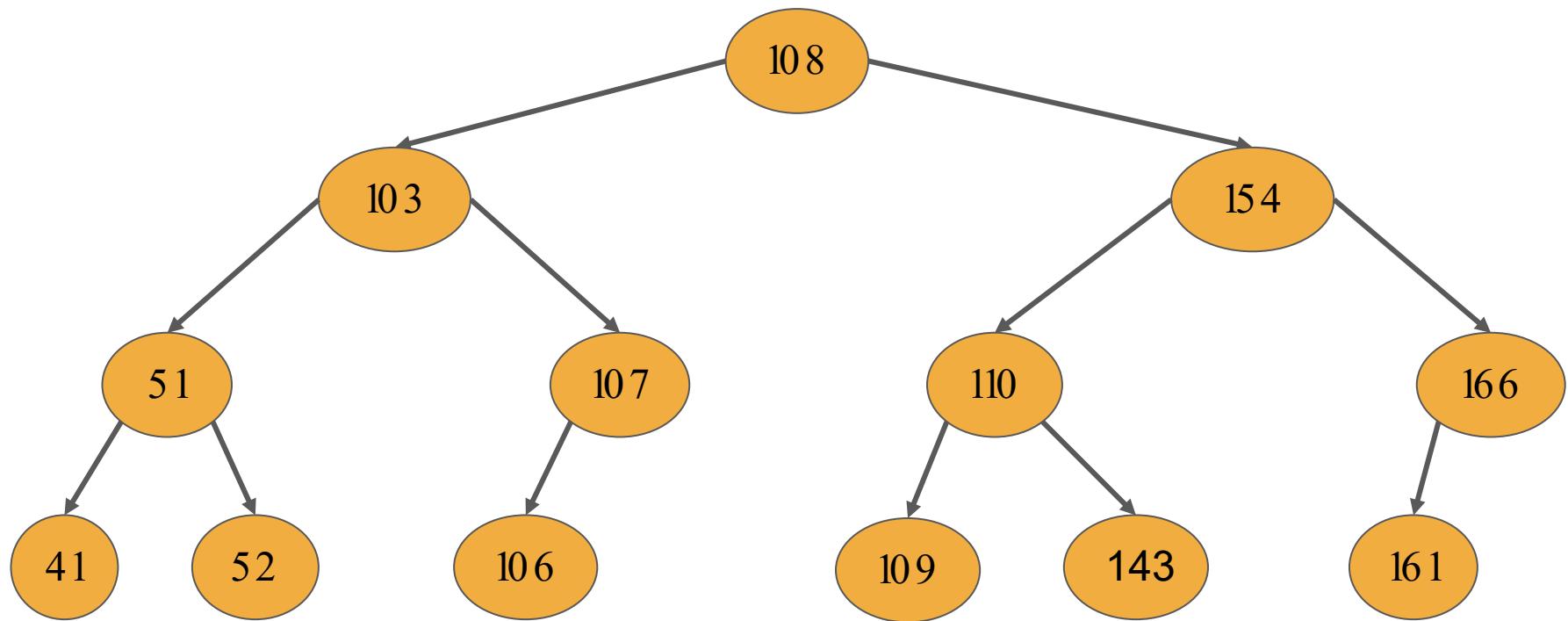
We found **106** so we're done!

# Let's code it!

`contains()`

106 < 108  
106 > 103  
106 < 107

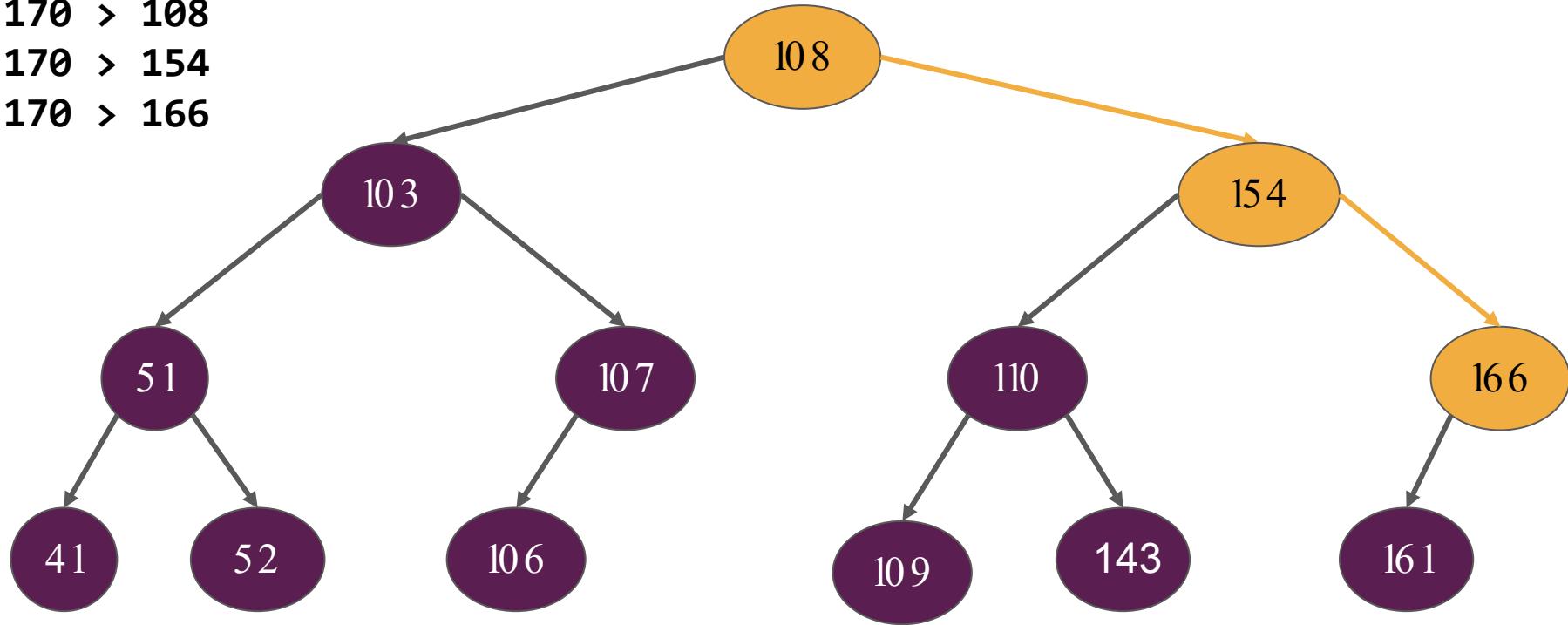




# Let's code it!

**add()**

170 > 108  
170 > 154  
170 > 166



Right child is **nullptr**

# OurSet summary

- Our tree utility functions (**inorderPrint**, **freeTree**) showed up as private member functions/helpers!
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# OurSet summary

- Our tree utility functions (**inorderPrint**, **freeTree**) showed up as private member functions/helpers!
  - In-order traversal prints our elements in the correctly sorted order!
- Using a BST allowed us to take advantage of recursion to traverse our data and get an  $O(\log n)$  runtime for our methods.
- Rewiring trees can be complicated!
  - Make sure to consider when nodes need to be passed by reference.
  - Check out the remove method after class if you're interested in seeing an example of tree rewiring (you won't be required to do anything this complex with tree rewiring).

You can play around with a simulation of BSTs to see traversal, lookup, insertion, and deletion, and try balanced and unbalanced trees at this website:

<https://visualgo.net/bn/bst>

# What's next?

# Roadmap

## Object-Oriented Programming

### C++ basics

User/client

vectors + grids

stacks + queues

sets + maps

Core Tools

testing

algorithmic analysis

recursive problem-solving

Diagnostic

Life after CS106

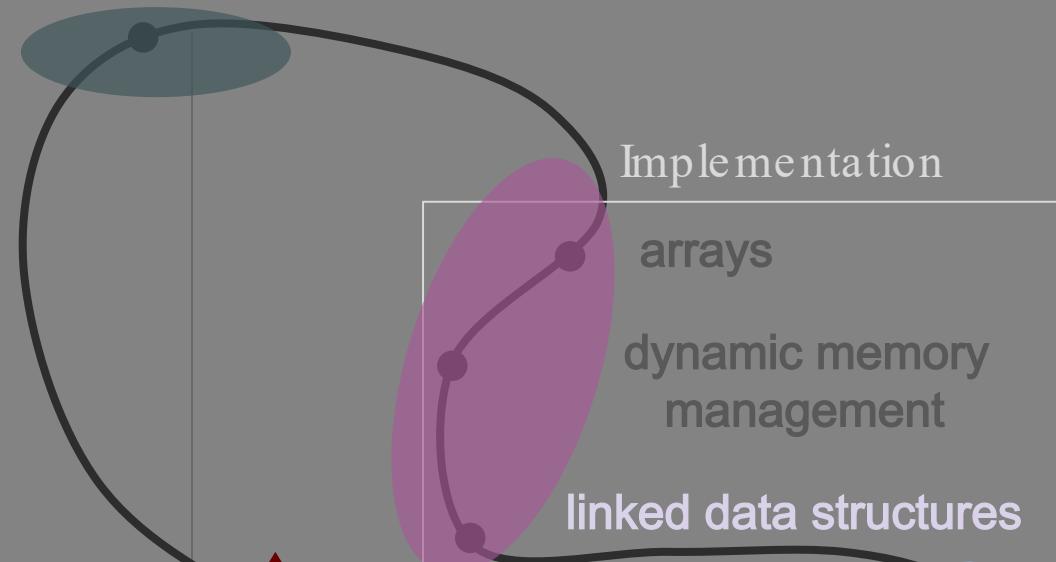
Implementation

arrays

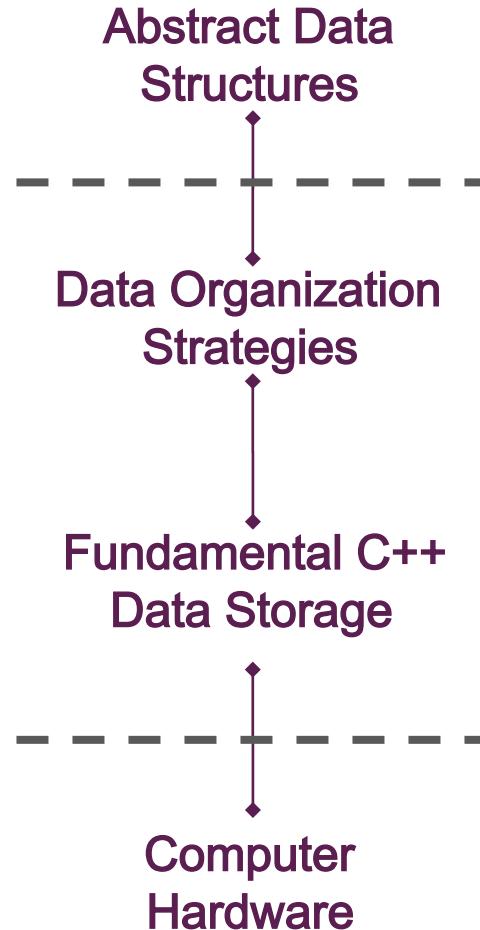
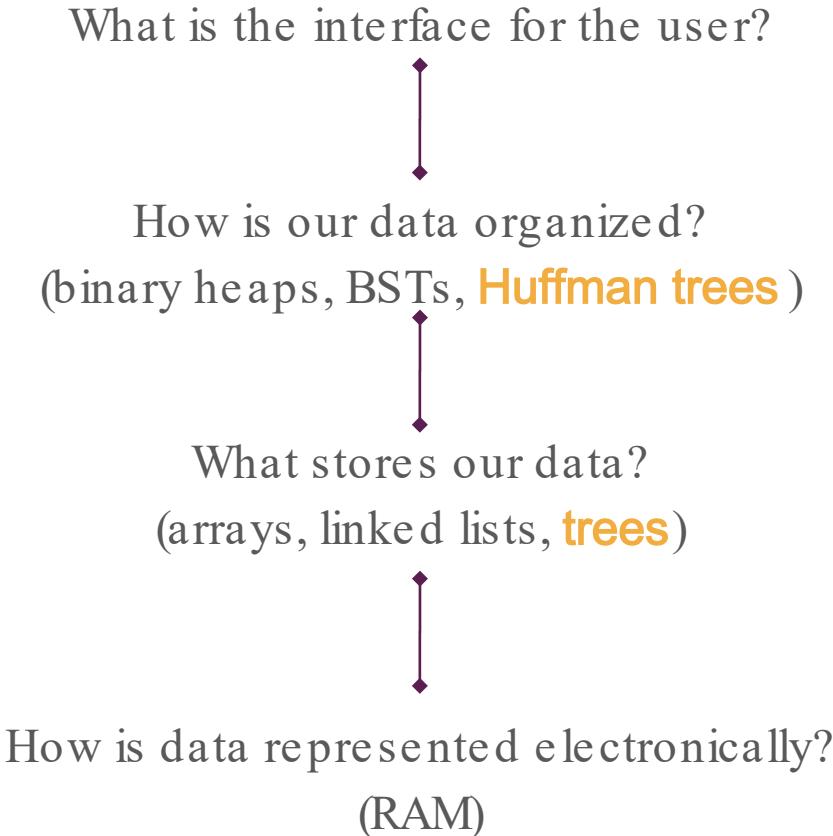
dynamic memory management

linked data structures

real-world algorithms



# Levels of abstraction



# Huffman coding



**RAW vs JPEG**

