Poll results

As of 1am this morning...

![Poll Results Diagram]
I have a bunch of practice problems.

Y’all vote on topics and we’ll do them.

I can also answer particular questions about the material.

Topics I have problems for:
- Grab-bag (multiple choice, etc)
- Hashing
- Red-Black Trees
- Ford-Fulkerson
- Dynamic Programming
- Greedy algorithms
- Divide and conquer
- Randomized algs

We ended up talking about these in class.
Multiple choice warmup!

For each of the following quantities, identify all of the options that correctly describe the quantity.

(a) The function $f(n)$, where $f(n) = n \log(n)$. \hspace{1em} (A), (C)

(b) $T(n)$ given by $T(n) = T(n/4) + \Theta(n^2)$ with $T(n) = 1$ for all $n \leq 8$. \hspace{1em} (A)(B)(C)

(c) $T(n)$ which is the running time of the following algorithm:

```python
mysteryAlg( n ):
    if n < 3:
        return 1
    return mysteryAlg( n/2 ) + mysteryAlg( (n/2) + 1 )
```

where above all division is integer division (so $a/b$ means $\lfloor a/b \rfloor$). \hspace{1em} (A), (C), (D)

(A) $O(n^2)$ \hspace{1em} (B) $\Theta(n^2)$ \hspace{1em} (C) $\Omega(n)$ \hspace{1em} (D) $O(n)$ \hspace{1em} (E) $O(\log^2(n))$. 
Let $G = (V, E)$ be an undirected weighted graph, and let $T$ be a minimum spanning tree in $G$. Decide whether the following statements must be true or may be false, and prove it!

(a) For any pair of distinct vertices $s, t \in V$, there is a unique path from $s$ to $t$ in $T$.

(b) For any pair of distinct vertices $s, t \in V$, the cost of a path between $s$ and $t$ in $T$ is minimal among all paths from $s$ to $t$ in $G$. 
Let $U$ be a universe of size $m$, where $m$ is a prime, and consider the following two hash families which hash $U$ into $n$ buckets, where $n$ is much smaller than $m$.

- First, consider $H_1$, which is the set of all functions from $U$ to $\{1, \ldots, n\}$:
  \[ H_1 = \{h \mid h : U \to \{1, \ldots, n\}\} \]

- Second, let $p = m$ (so $p$ is prime since we assumed $m$ to be prime), and choose $H_2$ to be
  \[ H_2 = \{h_{a,b} \mid a \in \{1, \ldots, p-1\}, b \in \{0, \ldots, p-1\}\}, \]
  where $h_{a,b}(x) = (ax + b \mod p) \mod n$.

You want to implement a hash table using one of these two families. Why would you choose $H_2$ over $H_1$? **Choose the best answer.**

(A) $H_1$ isn't a universal hash family.
(B) Storing an element of $H_1$ takes a lot of space.
(C) Storing all of $H_1$ takes a lot of space.
Shortest Paths

- When might you prefer breadth-first search to Dijkstra’s algorithm?
  
  **If the graph is unweighted**

- When might you prefer Floyd-Warshall to Bellman-Ford?
  
  **If you want shortest paths between all pairs of vertices.**

- When might you prefer Bellman-Ford to Dijkstra’s algorithm?
  
  **If there are negative edge weights**
Suppose that $b_1, \ldots, b_n$ are $n$ distinct integers in a uniformly random order. Consider the following algorithm:

```python
findMax(b_1, \ldots, b_n):
    currentMax = -Infinity
    for i = 1, \ldots, n:
        if b_i > currentMax:
            currentMax = b_i
    return currentMax
```

What is the expected number of times that `currentMax` is updated? (Asymptotic notation is fine).

$$
\mathbb{E}\left\{ \# \text{times currentMax updated} \right\} = \mathbb{E}\left\{ \sum_{i=1}^{n} \mathbb{I}\{b_i > b_1, \ldots, b_{i-1}\} \right\} = \sum_{i=1}^{n} \mathbb{P}\{b_i > b_1, \ldots, b_{i-1}\} = \sum_{i=1}^{n} \frac{1}{i} = \Theta(\log(n))
$$

Since $b_1, \ldots, b_i$ are uniformly random, the probability that any one is largest is $\frac{1}{i}$. 

$\Theta(\log(n))$
Consider the following flow on a graph. The notation $x/y$ means that an edge has flow $x$ out of capacity $y$.

Draw the residual graph for this flow.

Find an augmenting path in the residual graph and use it to increase the flow.

Find a minimum cut and prove (not by exhaustion) that it is a minimum cut.
Dynamic Programming!

- Suppose that roads in a city are laid out in an \( n \times n \) grid, but some of the roads are obstructed.
- For example, for \( n = 3 \), the city may look like this:

Define \( M[i,j] = \# \text{ paths from } (0,0) \text{ to } (i,j) \).

\[
M[i,j] = \begin{cases} 
\sum_{k=0}^{i-1} M[k,j] + \sum_{k=0}^{j-1} M[i,k] 
\end{cases} 
\]

\( \text{where } M[0,0] = 1 \).

- Design a DP algorithm to solve this problem.

where we have only drawn the roads that are not blocked. You want to count the number of ways to get from \((0,0)\) to \((n-1, n-1)\), using paths that only go up and to the right. In the example above, the number of paths is 3.
Given an array $A$ of length $n$, we say that an array $B$ is a *circular shift* of $A$ if there is an integer $k$ between 1 and $n$ (inclusive) so that


where $+$ denotes concatenation.

For example, if $A = [2, 5, 6, 8, 9]$, then $B = [6, 8, 9, 2, 5]$ is a circular shift of $A$ (with $k = 2$). The sorted array $A$ itself is also a circular shift of $A$ (with $k = 1$).

Design a $O(\log(n))$-time algorithm that takes as input an array $B$ which is a circular shift of a sorted array which contains distinct positive integers, and returns the value of the largest element in $B$. For example, give $B$ as above, your algorithm should return 9.

Solution on next page.
def findMax(B):
    n ← len(B)
    if B[0] ≤ B[n-1]:   \ case 1
        return B[n-1]
    mid = \[n/2\] + 1
    if B[mid] > B[0]:   \ case 2
        return findMax( B[mid:n] )
    if B[mid] < B[0]:   \ case 3
        return findMax( B[0:mid] )

Idea:

- In CASE 1, the situation looks like

So we return B[n-1]

- In CASE 2, it looks like

So the max is on the right side and we recurse on B[mid:

- In CASE 3, it looks like

So the max is on the left side and we recurse on B[mid+1]


Greedy Algorithms!

There are $n$ final exams on Dec. 13 at Stanford; exam $i$ is scheduled to begin at time $a_i$ and end at time $b_i$. Two exams which overlap cannot be administered in the same classroom; two exams $i$ and $j$ are defined to be overlapping if $[a_i, b_i] \cap [a_j, b_j] \neq \emptyset$ (including if $b_i = a_j$, so one starts exactly at the time that the other ends). Design an algorithm which solves the following problem.

- **Input:** Arrays $A$ and $B$ of length $n$ so that $A[i] = a_i$ and $B[i] = b_i$.
- **Output:** The smallest number of classrooms necessary to schedule all of the exams, and an optimal assignment of exams to classrooms.
- **Running time:** $O(n \log(n) + nk)$, where $k$ is the minimum number of classrooms needed.
- **For example:** Suppose there are three exams, with start and finish times as given below:

<table>
<thead>
<tr>
<th>$i$</th>
<th>1</th>
<th>2</th>
<th>3</th>
</tr>
</thead>
<tbody>
<tr>
<td>$a_i$</td>
<td>12pm</td>
<td>4pm</td>
<td>2pm</td>
</tr>
<tr>
<td>$b_i$</td>
<td>3pm</td>
<td>6pm</td>
<td>5pm</td>
</tr>
</tbody>
</table>

Then the exams can be scheduled in two rooms; Exam 1 and Exam 2 can be scheduled in Room 1 and Exam 3 can be scheduled in Room 2.
def scheduleRooms(A, B):
    n ← len(A)
    C = \{ (A[i], i) \mid i = 0, \ldots, n - 1 \}
    C = \text{sort in increasing order by start time.}
    rooms = [ ]
    endTimes = [ ]
    for i = 0, \ldots, n-1:
        C[i][1] > endTimes[i]:
            rooms[j].append(C[i][1])
            endTimes[j] = endTimes[j] + B[C[i][1]]
            break
    else:
        rooms.append(C[i][1])
        endTimes.append(B[C[i][1]])
    rooms

Correctness by induction

Inductive Hypothesis: After adding the i-th exam, there is an optimal schedule that extends the current solution.

Base Case: After adding 0 exams, there is an optimal schedule extending times.

Inductive Step: Suppose the inductive hypothesis holds for i-1, and let S be the optimal schedule that extends it.

If S puts exam i where we would put it (say, room r), then we are done, so suppose that S puts exam i in room r'.

Let j > i be the next exam scheduled in room r.

Then o_j >= o_i, since o_i had the smallest start time of all exams not yet picked.

So consider the schedule S' where we swap the rest of room r' with the rest of room r:

Room 1:
Room r:
Room r':

This is still a valid schedule, and uses the same number of rooms as S, so it's also optimal. And it puts exam i in room r' we're done.

Conclusion: At the end of the algorithm, there's still an optimal solution extending the current one to the current one is optimal.
Universal Hash Families

Definition: A hash family $\mathcal{H}$ (mapping $\mathcal{U}$ into $n$ buckets) is **2-universal** if for all $x \neq y \in \mathcal{U}$ and for all $a, b \in \{1, \ldots, n\}$,

$$\mathbb{P}((h(x), h(y)) = (a, b)) = \frac{1}{n^2}.$$ 

(a) Show that if $\mathcal{H}$ is 2-universal, then it is universal.

(b) Show that the converse is not true. That is, there is a universal family that’s not 2-universal.

(a) Suppose that $\mathcal{H}$ is 2-universal. Then $\forall x \neq y \in \mathcal{U}$,

$$\mathbb{P}\{h(x) = h(y)\} = \sum_{t \in \{1, \ldots, n\}} \mathbb{P}\{(h(x), h(y)) = (tt)\}$$

$$= \sum_{t \in \{1, \ldots, n\}} \frac{1}{n^2} = \frac{1}{n}.$$ 

So by definition $\mathcal{H}$ is universal.

(b) Consider: $\mathcal{U} = \{x, y\}$, $\mathcal{H} = \{h_1, h_2\}$, where:

\[
\begin{array}{c|cc}
 & x & y \\
\hline
h_1 & 0 & 0 \\
h_2 & 1 & 0 \\
\end{array}
\]

Then $\mathbb{P}\{h(x) = h(y)\} = \frac{1}{2}$. But $\mathbb{P}\{h(x) = 0, h(y) = 1\} = \frac{1}{2}$, not $\frac{1}{4}$. 


More universal hash families

Say that \( \mathcal{H} \) is a universal hash family, containing functions \( h : \mathcal{U} \rightarrow \{1, \ldots, n\} \). Consider the following game.

- You choose \( h \in \mathcal{H} \) uniformly at random and keep it secret.
- A bad guy chooses \( x \in \mathcal{U} \), and asks you for \( h(x) \). (You give it to them).
- The bad guy chooses \( y \in \mathcal{U} \setminus \{x\} \), and tries to get \( h(y) = h(x) \).
- If \( h(x) = h(y) \), the bad guy wins. Otherwise, you win.

One of the following two is true.

1. There is a universal hash family \( \mathcal{H} \) so that the bad guy wins with probability 1.
2. For any universal hash family \( \mathcal{H} \), the probability that the bad guy wins is at most \( 1/n \).

Which is true and why?

To see this, consider the hash family \( \mathcal{H} = \{h_1, h_2\} \) over \( \mathcal{U} = \{x, y, z\} \) with \( n = 2 \).

\[
\begin{array}{ccc}
\mathcal{H} & x & y \\
\mathcal{H} & 1 & 0 & 1 \\
\end{array}
\]

This is universal since
\[
\begin{align*}
\Pr[h(x) = h(y)] & = 1/2 \\
\Pr[h(x) = h(z)] & = 1/2 \\
\Pr[h(y) = h(z)] & = 0 < 1/4 .
\end{align*}
\]

Bad guy's algorithm:
- If \( h(x) = 1 \), choose \( z \).
- If \( h(x) = 0 \), choose \( y \).
Red-Black Trees

Which of the following can be colored as a red-black tree? Either give a coloring or explain why not.