

Outline

- Privacy
- Collaborative Game Theory
- Clustering

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What is privacy?

- one of society's most vital concerns
- central for e-commerce
- arguably the most crucial and far-reaching current challenge and mission of CS
- least understood scientifically (e.g., is it rational?)
- see, e.g., www.sims.berkeley.edu/~hal, [~/pam](http://www.sims.berkeley.edu/~pam), [Stanford Law Review, June 2000]

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some thoughts on privacy

- also an economic problem
- surrendering private information is either good or bad for you
- example: privacy vs. search costs in computer purchasing

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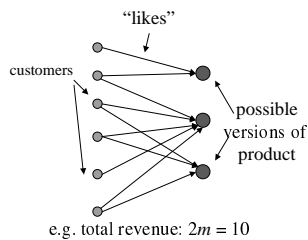
thoughts on privacy (cont.)

- *personal information is intellectual property controlled by others, often bearing negative royalty*
- selling mailing lists vs. selling aggregate information: false dilemma
- Proposal: *Take into account the individual's utility when using personal data for decision-making*

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e.g., marketing survey



- company's utility is proportional to the majority
- customer's utility is 1 if in the majority
- *how should all participants be compensated?*

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Collaborative Game Theory

- How should A, B, C split the loot (=20)?
 - We are given what each subset can achieve by itself as a function v from the powerset of $\{A, B, C\}$ to the reals
 - $v(\{\}) = 0$
- | Values of v | |
|---------------|----|
| • A: | 10 |
| • B: | 0 |
| • C: | 6 |
| • AB: | 14 |
| • BC: | 9 |
| • AC: | 16 |
| • ABC: | 20 |

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first idea (notion of “fairness”):
the core

A vector (x_1, x_2, \dots, x_n) with $\sum_i x_i = v([n]) (= 20)$
is in the core if for all S we have
 $x[S] \geq v(S)$

In our example: A gets 11, B gets 3, C gets 6

Problem: Core is often empty (e.g., AB ✱ 15)

second idea: the Shapley value

$$x_i = E_{\pi}(v[\{j: \pi(j) \leq \pi(i)\}] - v[\{j: \pi(j) < \pi(i)\}])$$

(Meaning: Assume that the agents arrive at
random. Pay each one his/her contribution.
Average over all possible orders of arrival.)

Theorem [Shapley]: The Shapley value is the
only allocation that satisfies Shapley’s axioms.

In our example...

- | | |
|-------------------------------|---------------|
| • A gets: | Values of v |
| $10/3 + 14/6 + 10/6 +$ | • A: 10 |
| $11/3 = 11$ | • B: 0 |
| • B gets: | • C: 6 |
| $0/3 + 4/6 + 3/6 + 4/3 = 2.5$ | • AB: 14 |
| • C gets the rest = 6.5 | • BC: 9 |
| • NB: Split the cost of a | • AC: 16 |
| trip among hosts... | • ABC: 20 |

e.g., the UN security council

- 5 permanent, 10 non-permanent
- A resolution passes if voted by a majority of the 15, including all 5 P
- $v[S] = 1$ if $|S| > 7$ and S contains 1,2,3,4,5; otherwise 0
- What is the Shapley value (\sim power) of each P member? Of each NP member?

e.g., the UN security council

- What is the probability, when you are the 8th arrival, that all of 1,...,5 have arrived?
- Ans: $\text{Choose}(10,2)/\text{Choose}(15,7) \sim .7\%$
Permanent members: $\sim 18\%$

Therefore, P \oplus NP

third idea: bargaining set

fourth idea: nucleolus

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seventeenth idea: the von Neumann-Morgenstern solution

[Deng and P. 1990] complexity-theoretic
critique of solution concepts

Applying to the market survey problem

- Suppose largest *minority* is r
- An allocation is in the core as long as losers get 0, vendor gets $> 2r$, winners split an amount up to twice their victory margin
- (plus another technical condition saying that split must not be too skewed)

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market survey problem: Shapley value

- Suppose margin of victory is at least $\mu > 0\%$
- (realistic, close elections never happen in real life)
- Vendor gets $m(1 + \mu)$
- Winners get $1 + \mu$
- Losers get μ
- (and so, no compensation is necessary)

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e.g., recommendation system

- Each participant i knows a set of items B_i
- Each benefits 1 from every new item
- Core: empty, unless the sets are disjoint!
- Shapley value: For each item you know, you are owed an amount equal to $1 / (\# \text{people who know about it})$
--i.e., *novelty pays*

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e.g., collaborative filtering

- Each participant likes/dislikes a set of items (participant is a vector of 0, ± 1)
- The “similarity” of two agents is the inner product of their vectors
- There are k “well separated types” (vectors of ± 1), and each agent is a *random perturbation* and *random masking* of a type

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collaborative filtering (cont.)

- An agent gets advice on a 0 by asking the most similar other agent who has a ± 1 in that position
- Value of this advice is the product of the agent’s true value and the advice.
- How should agents be compensated (or charged) for their participation?

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collaborative filtering (result)

Theorem: An agent’s compensation (= value to the community) is an increasing function of how typical (close to his/her type) the agent is.

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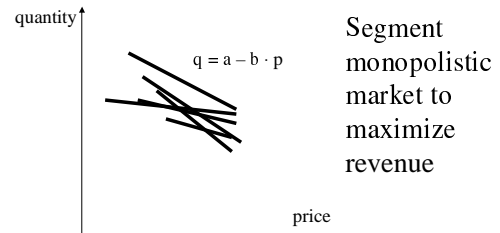
The economics of clustering

- The practice of clustering: Confusion, too many criteria and heuristics, no guidelines
- The theory of clustering: ditto!
- “It’s the economy, stupid!”
[Kleinberg, P., Raghavan STOC 98, JDKD 99]

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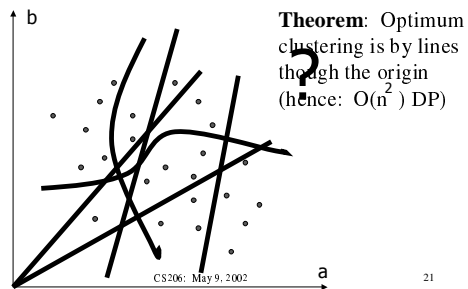
Example: market segmentation



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or, in the $a - b$ plane:



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So...

- Privacy has an interesting (and, I think, central) economic aspect
- Which gives rise to neat math/algorithmic problems
- Architectural problems wide open
- And clustering is a meaningful problem only in a well-defined economic context

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