

# CS234 Problem Session Solutions

Week 7: Feb 24

## 1) [CA Session] Useful Probability Bounds

In this problem, we will derive bounds to answer questions of the form: given a random variable  $Z$  with expectation  $\mathbb{E}[Z]$ , how likely is  $Z$  to be close to its expectation?

- (a) First, we will prove Markov's inequality: Let  $Z \geq 0$  be a non-negative random variable. Prove that for all  $t \geq 0$ ,

$$\mathbb{P}(Z \geq t) \leq \frac{\mathbb{E}[Z]}{t}$$

**Solution** First, notice that  $\mathbb{P}(Z \geq t) = \mathbb{E}[\mathbf{1}\{Z \geq t\}]$ , and that if  $Z \geq t$ , then it must be that  $\frac{Z}{t} \geq \mathbf{1}\{Z \geq t\}$ . Otherwise, if  $Z < t$ , then we have that  $\frac{Z}{t} \geq 0 = \mathbf{1}\{Z \geq t\}$ . Thus, we have that

$$\mathbb{P}(Z \geq t) = \mathbb{E}[\mathbf{1}\{Z \geq t\}] \leq \mathbb{E}\left[\frac{Z}{t}\right] = \frac{\mathbb{E}[Z]}{t}, \text{ as required.}$$

- (b) Next, we will prove Chebyshev's inequality. Let  $Z$  be any random variable with  $\text{Var}(Z) < \infty$ . Prove that for all  $t \geq 0$ ,

$$\mathbb{P}(Z \geq \mathbb{E}[Z] + t \text{ or } Z \leq \mathbb{E}[Z] - t) \leq \frac{\text{Var}(Z)}{t^2}$$

**Solution** This result follows from Markov's inequality. Notice that if  $Z \geq \mathbb{E}[Z] + t$ , then it is also true that  $(Z - \mathbb{E}[Z])^2 \geq t^2$ . Similarly, if  $Z \leq \mathbb{E}[Z] - t$ , then we have  $(Z - \mathbb{E}[Z])^2 \geq t^2$ . Hence, by Markov's inequality, we have that

$$\mathbb{P}(Z \geq \mathbb{E}[Z] + t \text{ or } Z \leq \mathbb{E}[Z] - t) = \mathbb{P}((Z - \mathbb{E}[Z])^2 \geq t^2) \leq \frac{\mathbb{E}[(Z - \mathbb{E}[Z])^2]}{t^2} = \frac{\text{Var}(Z)}{t^2}$$

- (c) It can be useful to derive tighter bounds through exponentially decreasing functions. Let us define the moment generating function for a random variable  $Z$  as

$$M_Z(\lambda) = \mathbb{E}[\exp(\lambda Z)]$$

We will now prove the Chernoff bound. Let  $Z$  be a random variable. Prove that for any  $t \geq 0$ ,

$$\mathbb{P}(Z \geq \mathbb{E}[Z] + t) \leq \min_{\lambda \geq 0} \mathbb{E}[e^{\lambda(Z - \mathbb{E}[Z])}] e^{-\lambda t} = \min_{\lambda \geq 0} M_{Z - \mathbb{E}[Z]}(\lambda) e^{-\lambda t}$$

**Solution** We will again prove this using Markov's inequality. For any  $\lambda > 0$ , we see that  $Z \geq \mathbb{E}[Z] + t$  if and only if  $e^{\lambda Z} \geq e^{\lambda(\mathbb{E}[Z]+t)}$ . Rearranging, we have  $e^{\lambda(Z-\mathbb{E}[Z])} \geq e^{\lambda t}$ . Now, we can apply Markov's inequality and see that

$$\mathbb{P}(Z - \mathbb{E}[Z] \geq t) = \mathbb{P}(e^{\lambda(Z-\mathbb{E}[Z])} \geq e^{\lambda t}) \leq \mathbb{E}[e^{\lambda(Z-\mathbb{E}[Z])}]e^{-\lambda t}$$

Notice that this bound certainly holds if  $\lambda = 0$ . Further, we have proven this for an arbitrary non-negative  $\lambda$ , so we can minimize the bound with respect to  $\lambda$  to achieve the tightest bound.

## 2) [Breakout Rooms] KL Divergence

The Kullback-Leibler (KL) divergence is defined as a measure of how different a probability distribution is from a second reference probability distribution. For discrete probability distributions  $P$  and  $Q$  defined over the same probability space  $X$ , the KL divergence is defined as

$$D_{KL}(P||Q) = \sum_{x \in X} P(x) \log\left(\frac{P(x)}{Q(x)}\right)$$

Show that the KL divergence is guaranteed to be non-negative.

**Solution** This can be proven in a few ways. We will prove this using Jensen's inequality. We will show that  $-D_{KL}(P||Q) \leq 0$ .

$$\begin{aligned} -D_{KL}(P||Q) &= -\sum_{x \in X} P(x) \log\left(\frac{P(x)}{Q(x)}\right) \\ &= \sum_{x \in X} P(x) \log\left(\frac{Q(x)}{P(x)}\right) \\ &\leq \log \sum_{x \in X} P(x) \frac{Q(x)}{P(x)} \text{ by Jensen's inequality since } \log \text{ is concave.} \\ &= \log \sum_{x \in X} Q(x) \\ &= \log(1) \\ &= 0 \end{aligned}$$

### 3) [Breakout Rooms] Probably Approximately Correct

Let  $A(\alpha, \beta)$  be a hypothetical reinforcement learning algorithm, parametrized in terms of  $\alpha$  and  $\beta$  such that for any  $\alpha > \beta > 1$ , it selects action  $a$  for state  $s$  satisfying  $|Q(s, a) - V^*(s)| \leq \frac{\beta}{\alpha}$  in all but  $N = \frac{|S||A|\alpha\beta}{1-\gamma}$  steps with probability at least  $1 - \frac{1}{\beta^2}$ .

Per the definition of *Probably Approximately Correct Reinforcement Learning*, express  $N$  as a function of  $|S|$ ,  $|A|$ ,  $\delta$ ,  $\epsilon$  and  $\gamma$ . What is the resulting  $N$ ? Is algorithm  $A$  probably approximately correct? Briefly justify.

**Solution** We want to achieve the bound that  $|Q(s, a) - V^*(s)| \leq \epsilon$  with probability  $1 - \delta$ . So let  $\frac{\beta}{\alpha} = \epsilon$  and  $1 - \frac{1}{\beta^2} = 1 - \delta$ , which gives  $\alpha = \frac{1}{\epsilon\sqrt{\delta}}$  and  $\beta = \frac{1}{\sqrt{\delta}}$ .

Substituting,  $N = \frac{|S||A|\alpha\beta}{1-\gamma} = \frac{|S||A|}{\epsilon\delta(1-\gamma)}$ .

Since  $N$  is a polynomial function of  $|S|$ ,  $|A|$ ,  $\frac{1}{\epsilon}$  and  $\frac{1}{\delta}$  and achieves the  $\epsilon, \delta$  bounds above, then  $A$  is PAC.