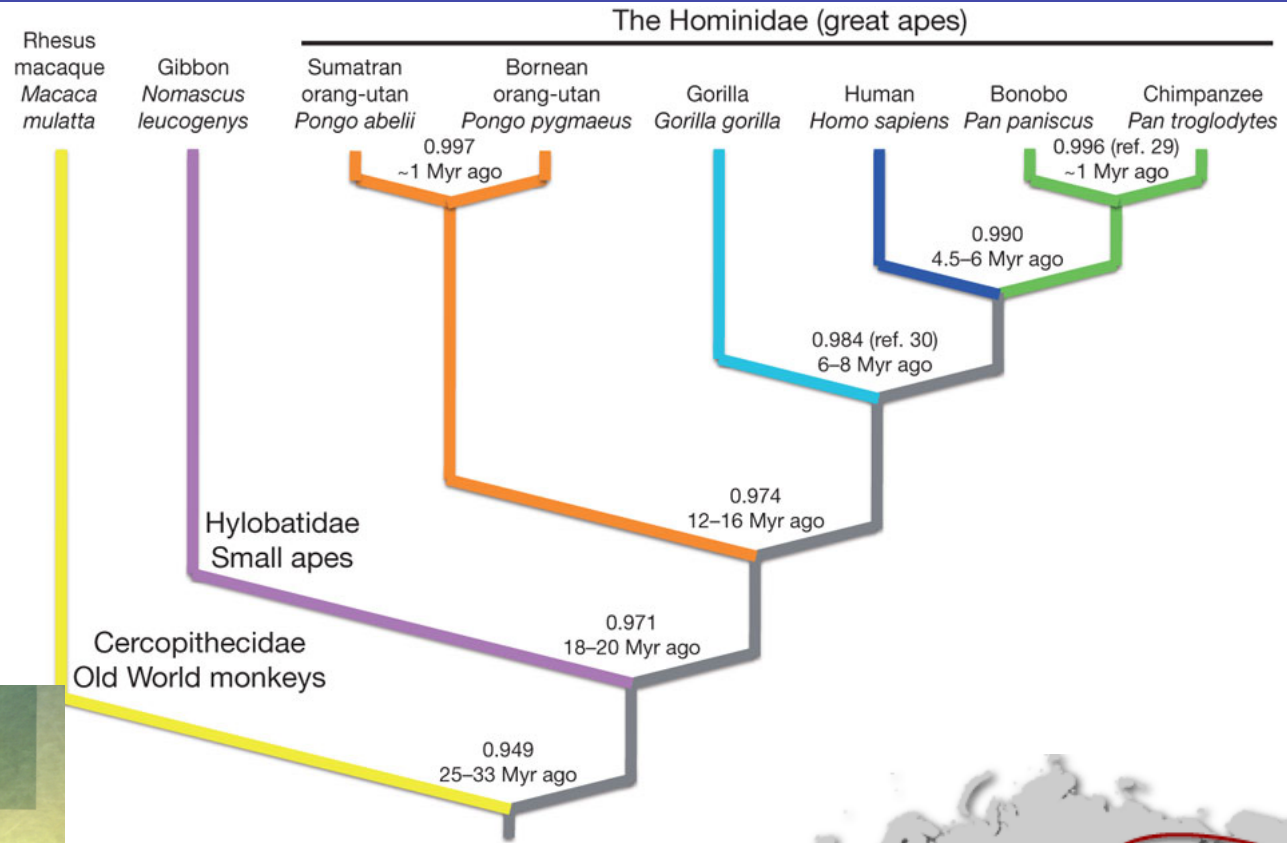


TGACTGAGGAGTTTACGGGAGCAAAGCGGCGTCATTGCTATTTCGTATCTGTTTAG  
CTGTGTTTACGGGAGCAAAGCGGCGTCATTGCTATTTCGTATCTGTTTAG

# Human Genome Diversity, Coalescence & Haplotypes



# The Hominid Lineage



**The Economist**

December 24th 2005 - January 8th 2006

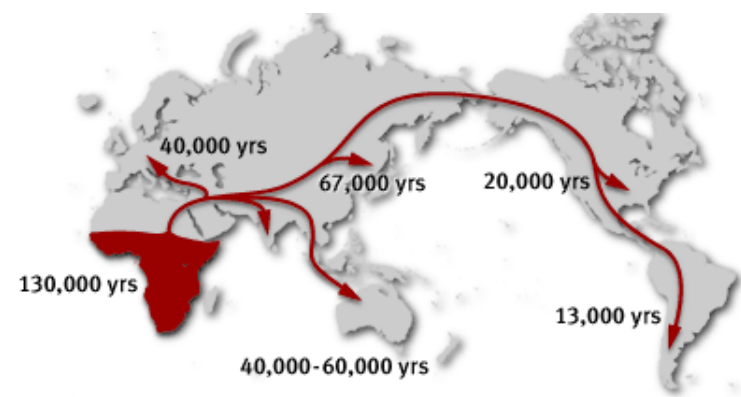
- Eavesdropping in America PAGE 33
- South Korea's cloning scandal PAGES 12 AND 103
- The dismal WTO deal PAGE 97
- Santa Claus, branding genius PAGE 90

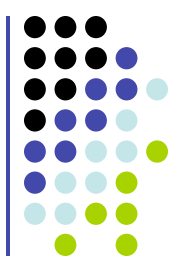
**Human Evolution**

What were our ancestors like?

Where did we evolve? Why big brains?

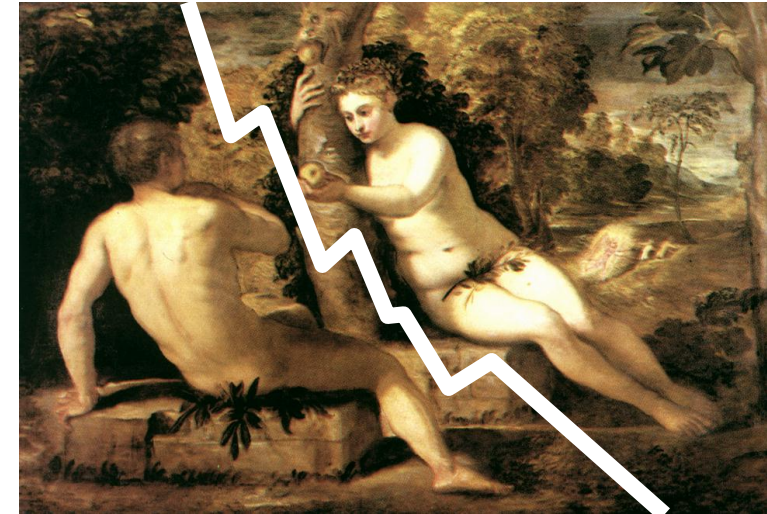
Relationships between populations?



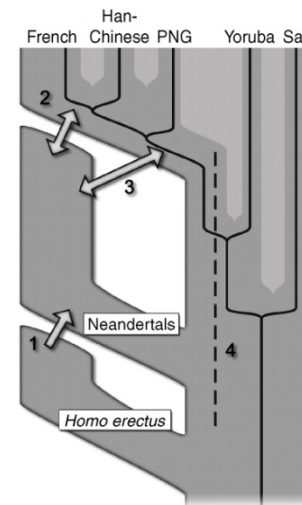


# Human population migrations

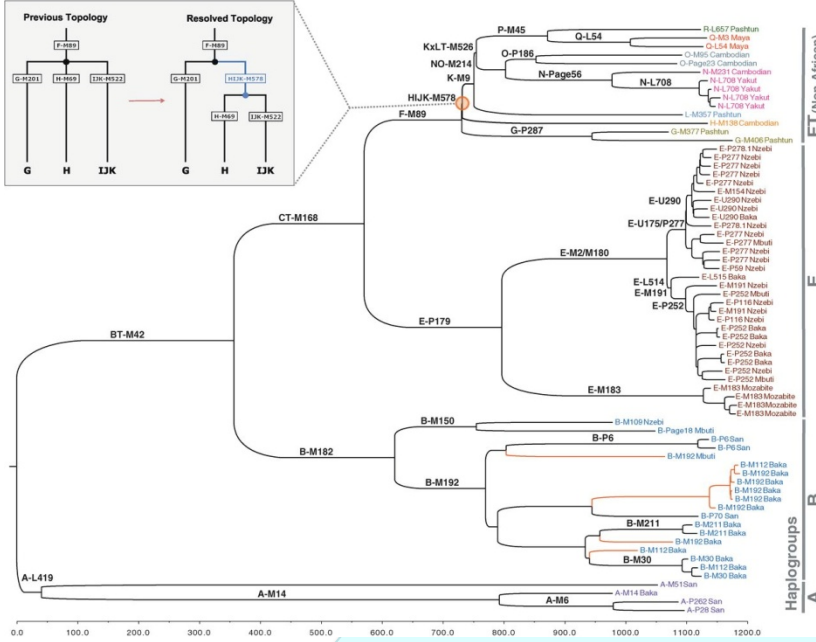
- Out of Africa, Replacement
  - Single mother of all humans (Eve)  
~99,000 – 150,000yr
  - Single father of all humans (Adam)  
~120,000 - 340,000yr
  - Humans out of Africa ~50000 years ago replaced others (e.g., Neandertals)



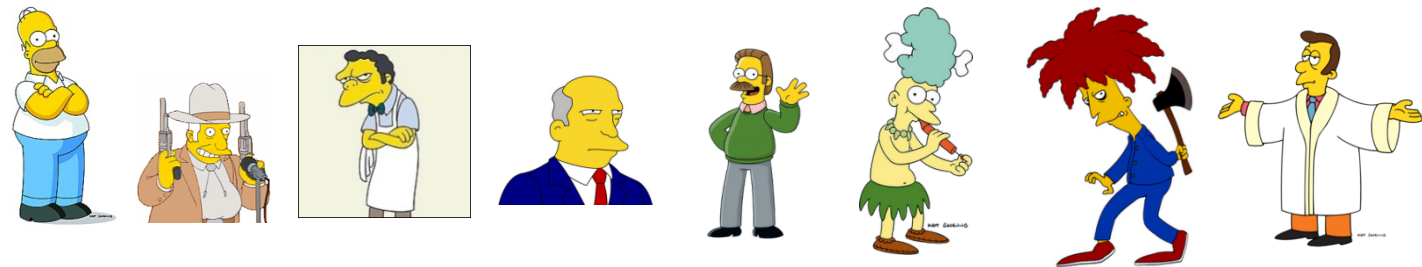
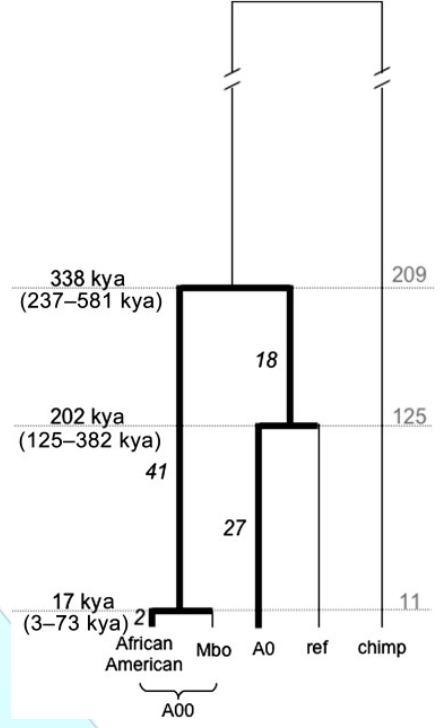
- Multiregional Evolution
  - Generally debunked, however,
  - ~5% of human genome in Europeans, Asians is Neanderthal, Denisova



# Coalescence



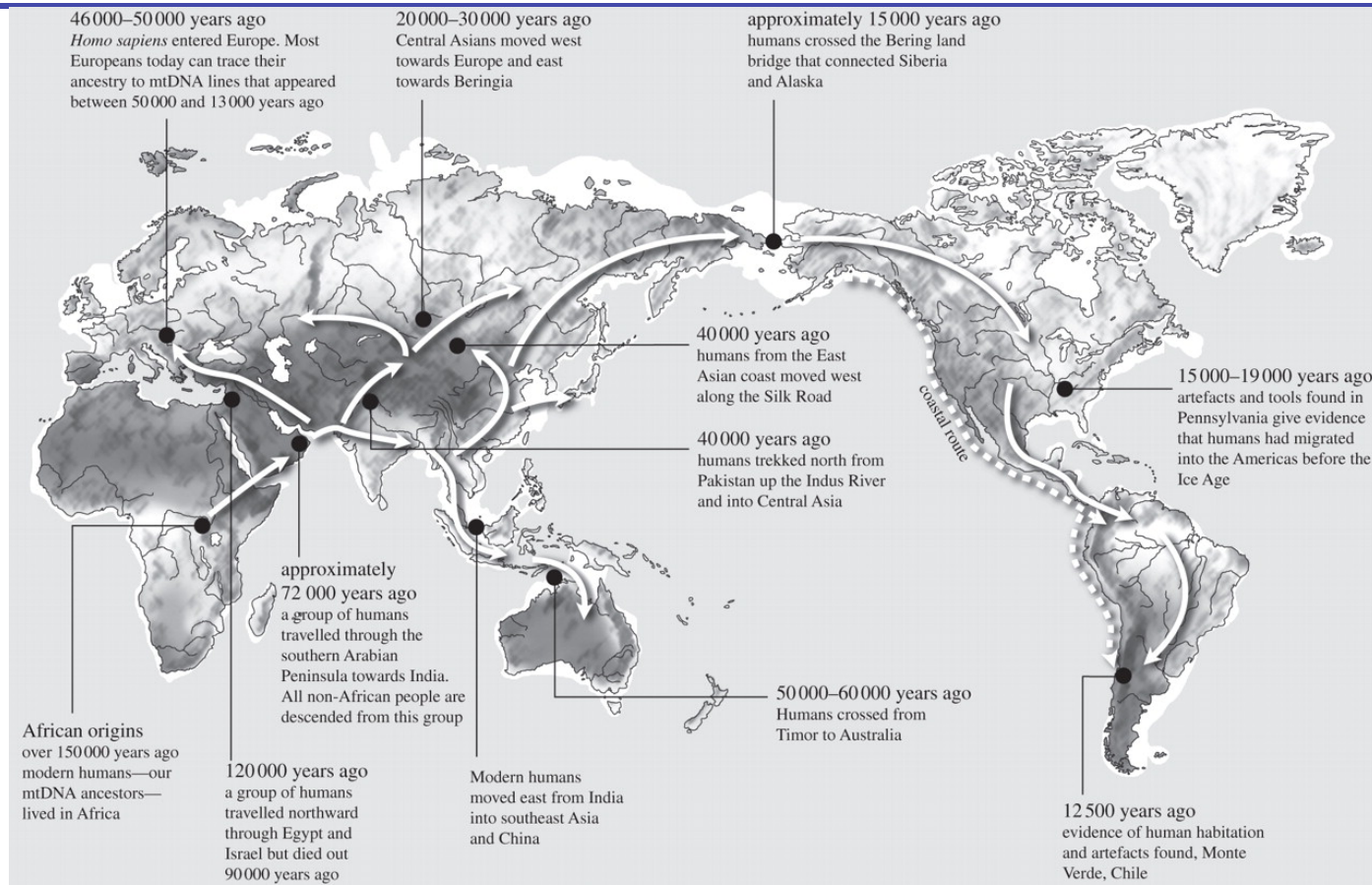
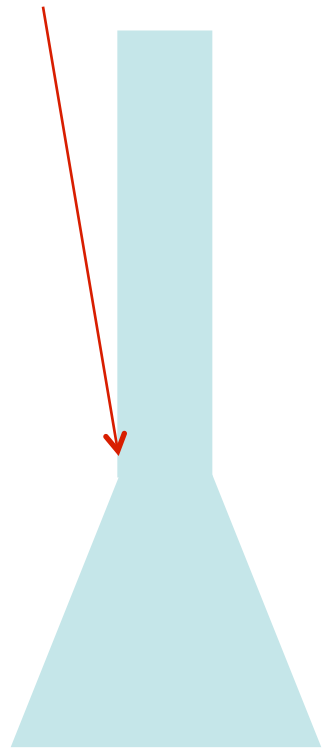
some coalescence





# Why humans are so similar

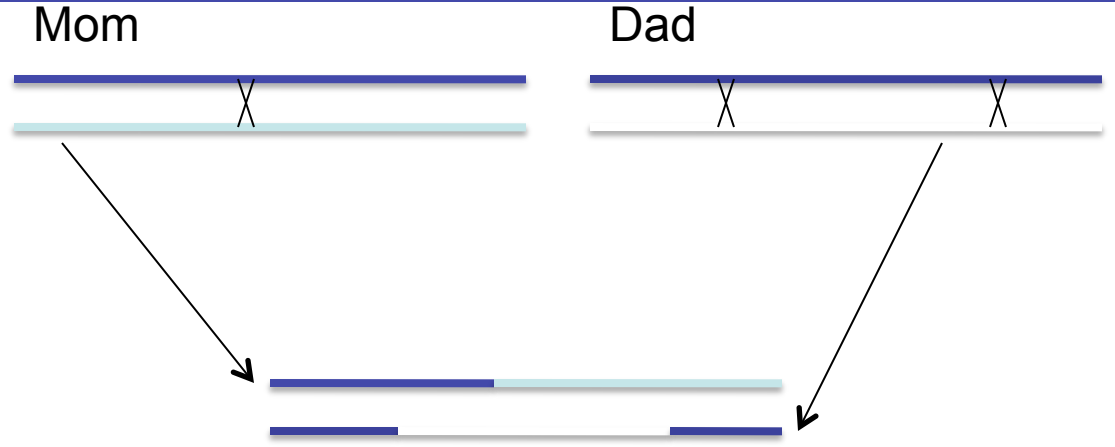
## Out of Africa





# Some Key Definitions

Mary:	AGCC	G/G	CG
John:	AGCC	G/G	CG
Josh:	AGCC	G/T	CG
Kate:	AGCC	G/G	CG
Pete:	AGCC	G/G	CG
Anne:	AGCC	G/G	CG
Mimi:	AGCC	G/G	CG
Mike:	AGCC	T/T	CG
Olga:	AGCC	T/G	CG
Tony:	AGCC	T/G	CG



**Alleles:** G, T

**Major Allele:** G

**Minor Allele:** T

**Heterozygosity:**  
 Prob[2 alleles picked at random with replacement are different]

$$2 * .75 * .25 = .375$$

$$H = 4Nu / (1 + 4Nu)$$

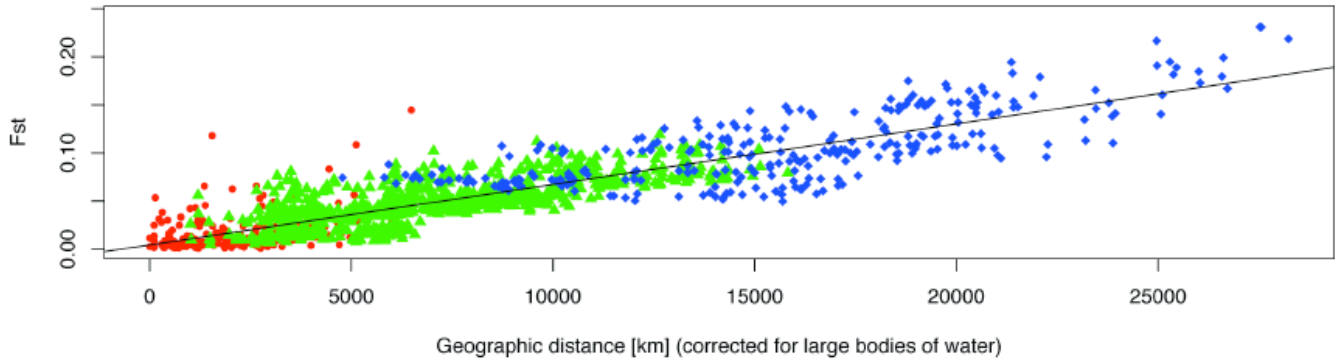
**Recombinations:**  
 At least 1/chromosome  
 On average ~1/100 Mb

**Linkage Disequilibrium:**  
 The degree of correlation between two SNP locations



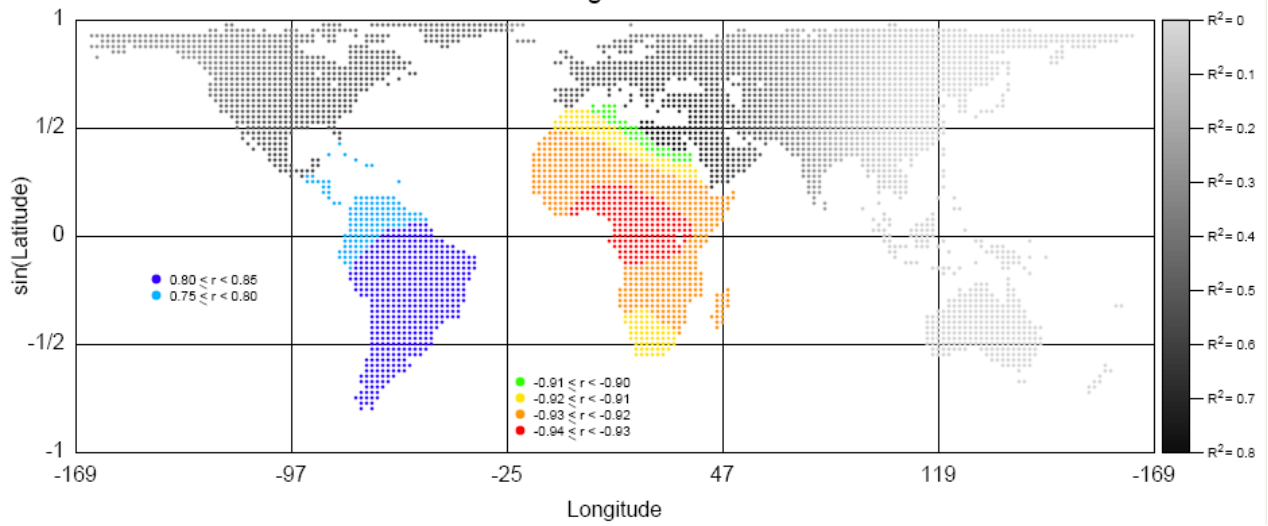
# The Fall in Heterozygosity

Figure 1B



$$F_{ST} = \frac{H - H_{POP}}{H}$$

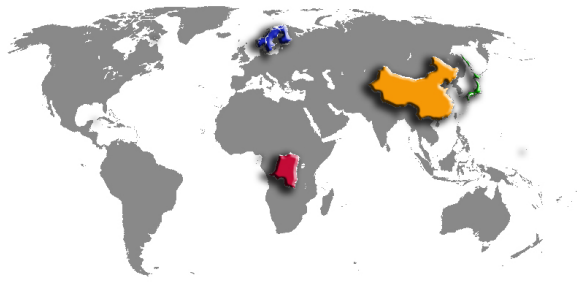
Figure 5



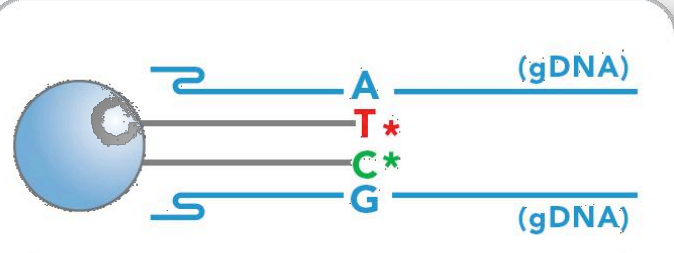


# The HapMap Project

ASW	African ancestry in Southwest USA	90
CEU	Northern and Western Europeans (Utah)	180
CHB	Han Chinese in Beijing, China	90
CHD	Chinese in Metropolitan Denver	100
GIH	Gujarati Indians in Houston, Texas	100
JPT	Japanese in Tokyo, Japan	91
LWK	Luhya in Webuye, Kenya	100
MXL	Mexican ancestry in Los Angeles	90
MKK	Maasai in Kinyawa, Kenya	180
TSI	Toscani in Italia	100
YRI	Yoruba in Ibadan, Nigeria	100



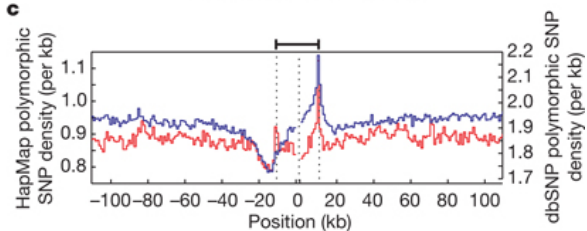
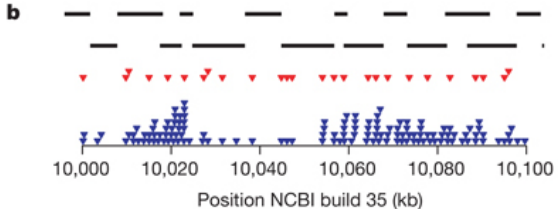
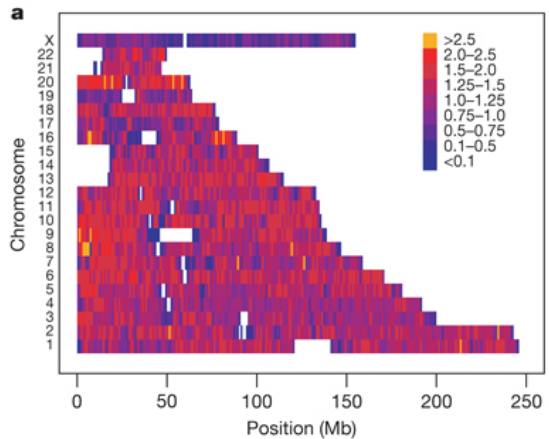
Genotyping:  
Probe a limited number (~1M) of known highly variable positions of the human genome



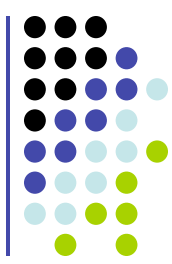
**7** Extend/Stain samples on BeadChip



**8** Image BeadChip





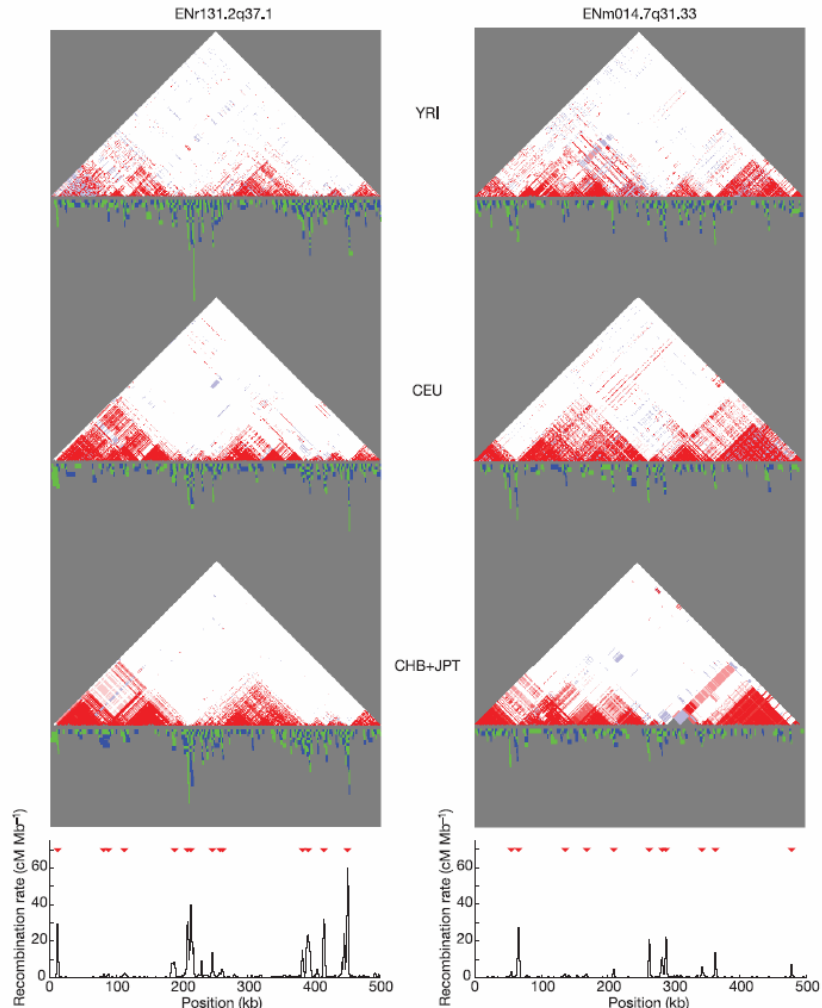


# Linkage Disequilibrium & Haplotype Blocks

Minor allele: A G  
 $p_A$   $p_G$

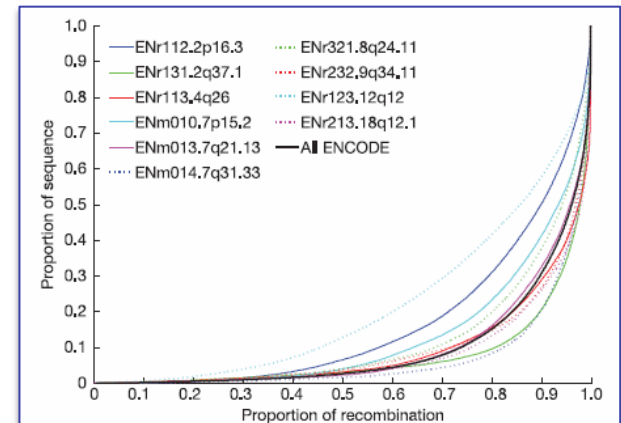
Linkage Disequilibrium (LD):

$$D = P(\text{A and G}) - p_A p_G$$



**Figure 8 | Comparison of linkage disequilibrium and recombination for two ENCODE regions.** For each region (ENr131.2q37.1 and ENm014.7q31.33),  $D'$  plots for the YRI, CEU and CHB+JPT analysis panels are shown: white,  $D' < 1$  and  $\text{LOD} < 2$ ; blue,  $D' = 1$  and  $\text{LOD} < 2$ ; pink,  $D' < 1$  and  $\text{LOD} \geq 2$ ; red,  $D' = 1$  and  $\text{LOD} \geq 2$ . Below each of these plots is shown the

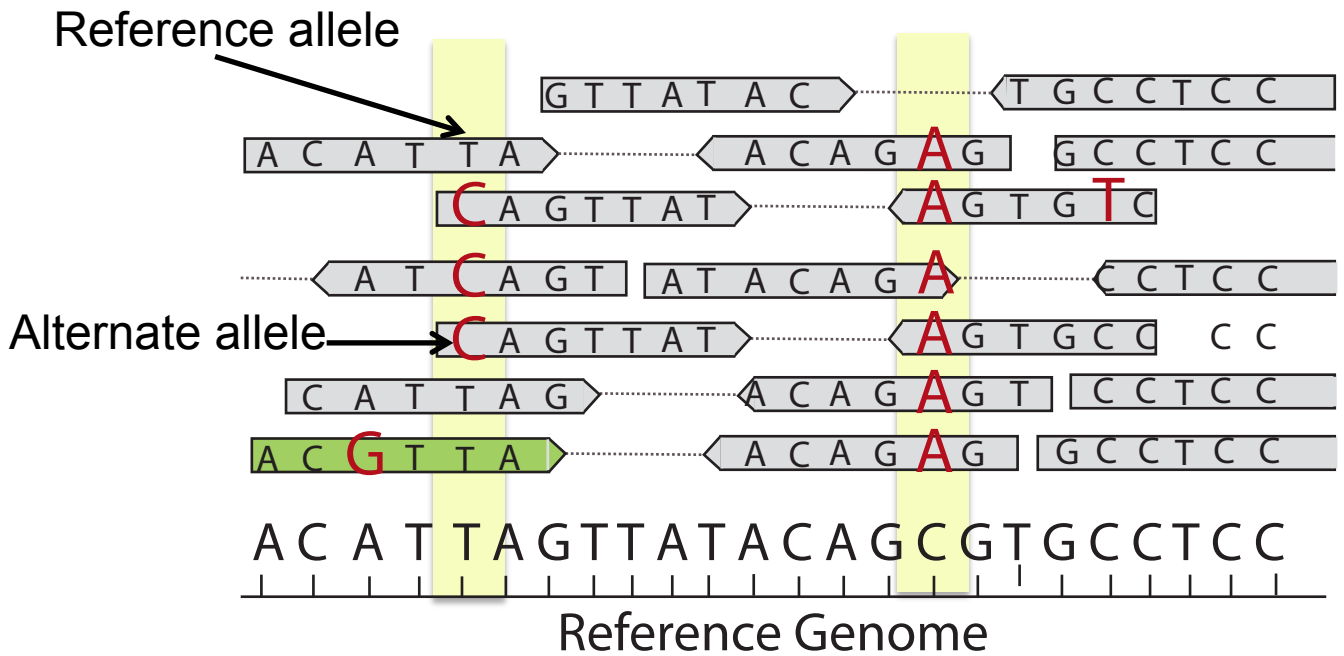
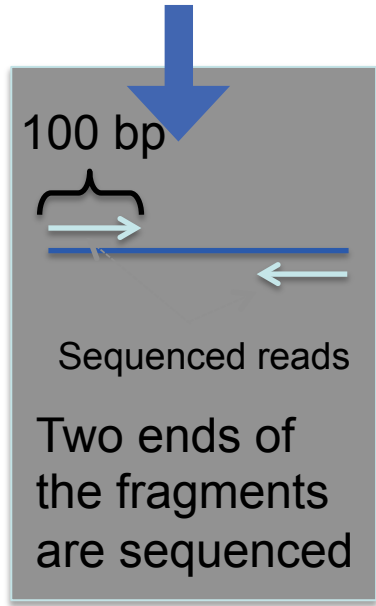
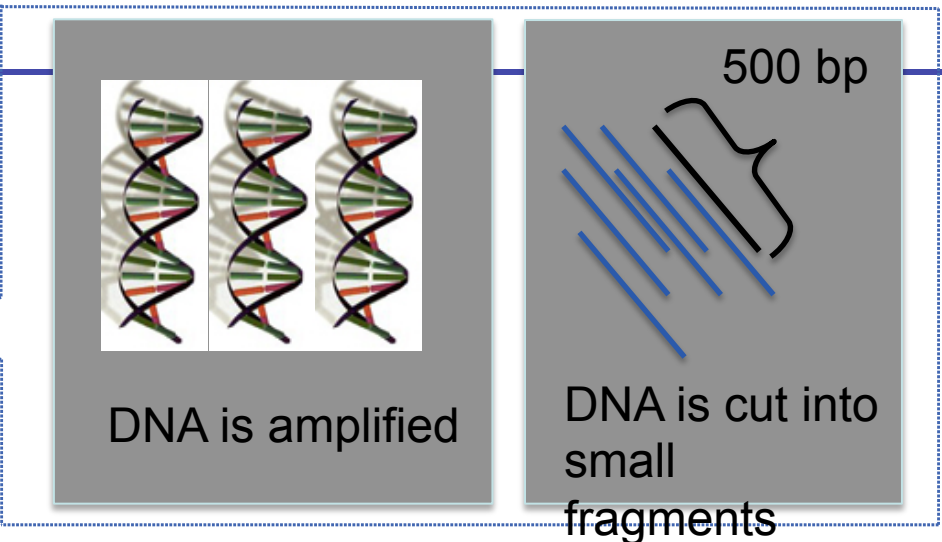
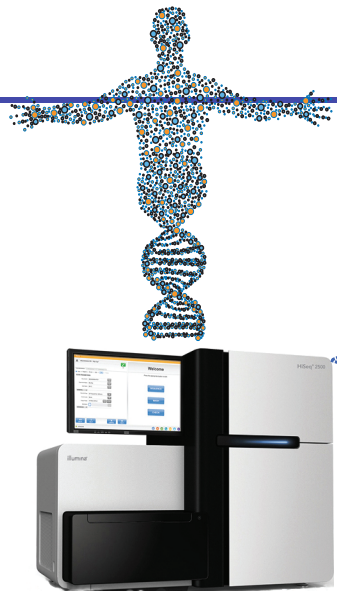
intervals where distinct obligate recombination events must have occurred (blue and green indicate adjacent intervals). Stacked intervals represent regions where there are multiple recombination events in the sample history. The bottom plot shows estimated recombination rates, with hotspots shown as red triangles<sup>46</sup>.



**Figure 9 | The distribution of recombination events over the ENCODE regions.** Proportion of sequence containing a given fraction of all recombination for the ten ENCODE regions (coloured lines) and combined across analysis panels (black line). For each line, SNP intervals are placed in decreasing order of estimated recombination rate<sup>46</sup>, combined across analysis panels, and the cumulative recombination fraction is plotted against the cumulative proportion of sequence. If recombination rates were constant, each line would lie exactly along the diagonal, and so lines further to the right reveal the fraction of regions where recombination is more strongly locally concentrated.

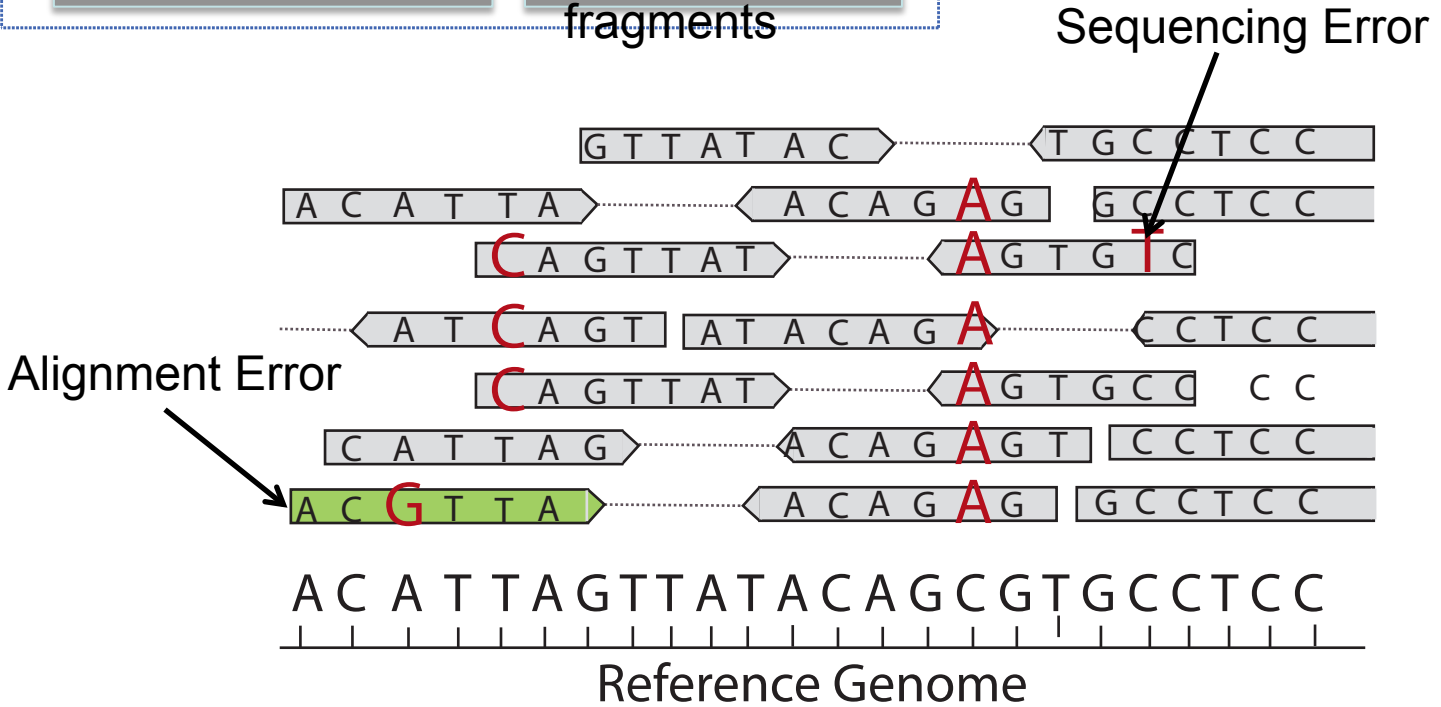
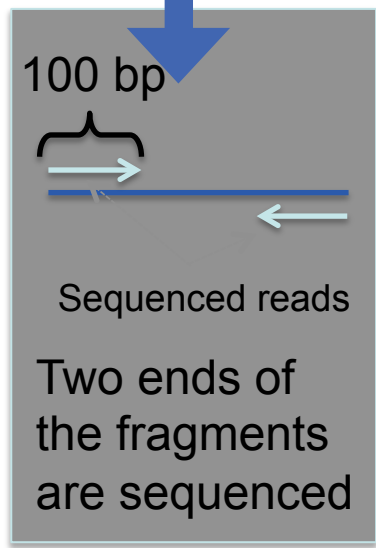
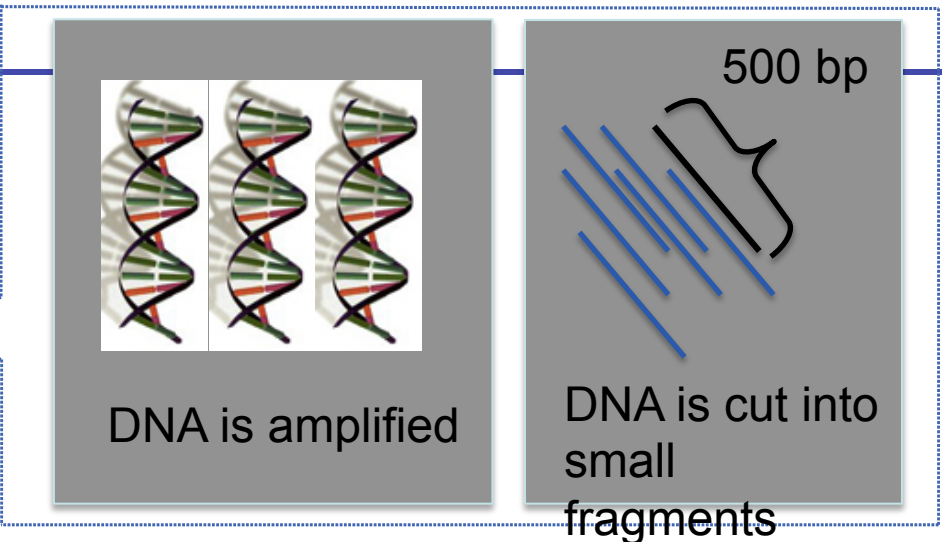
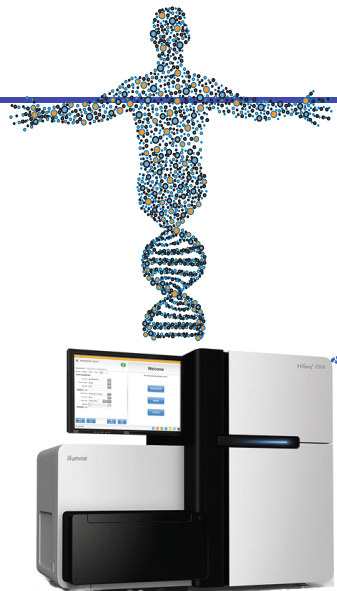


# Human Genome Resequencing





# Human Genome Resequencing



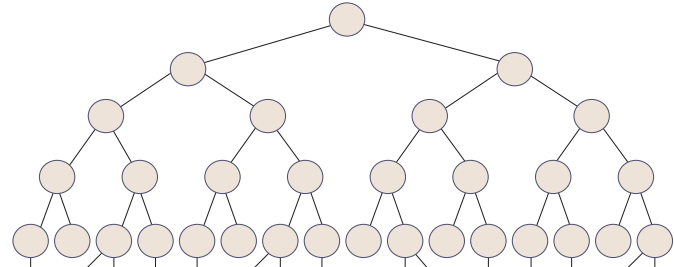


# Human Genome Resequencing

Types of SNVs :

## 1. Germline (SNPs)

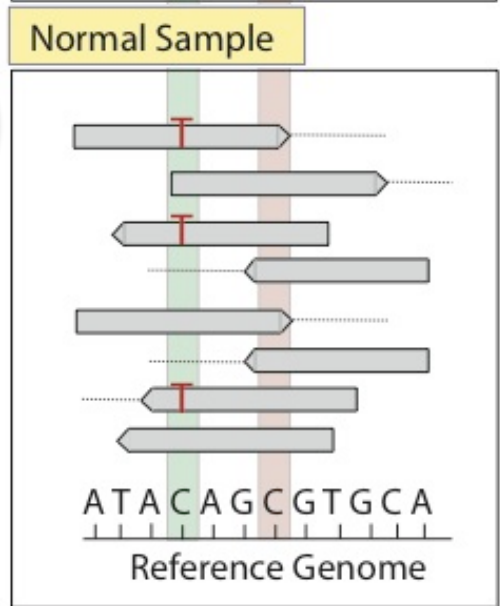
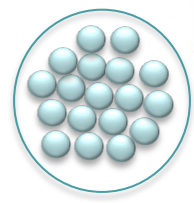
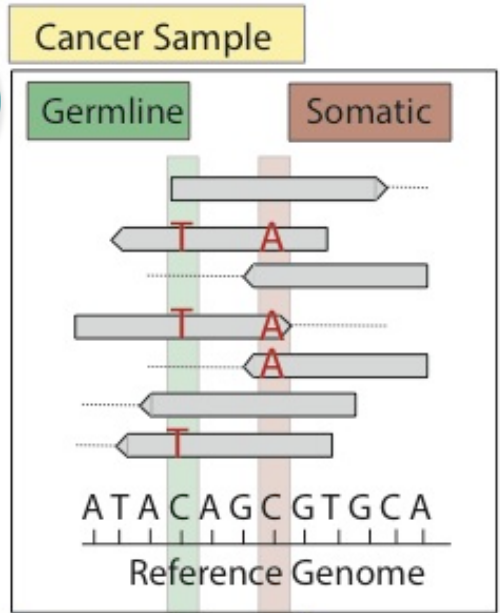
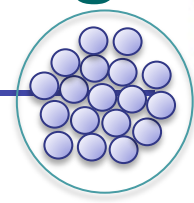
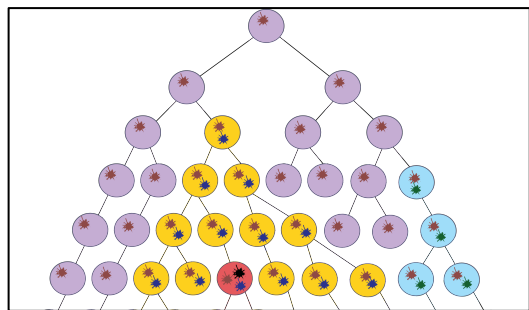
- Inherited
- All cells have it



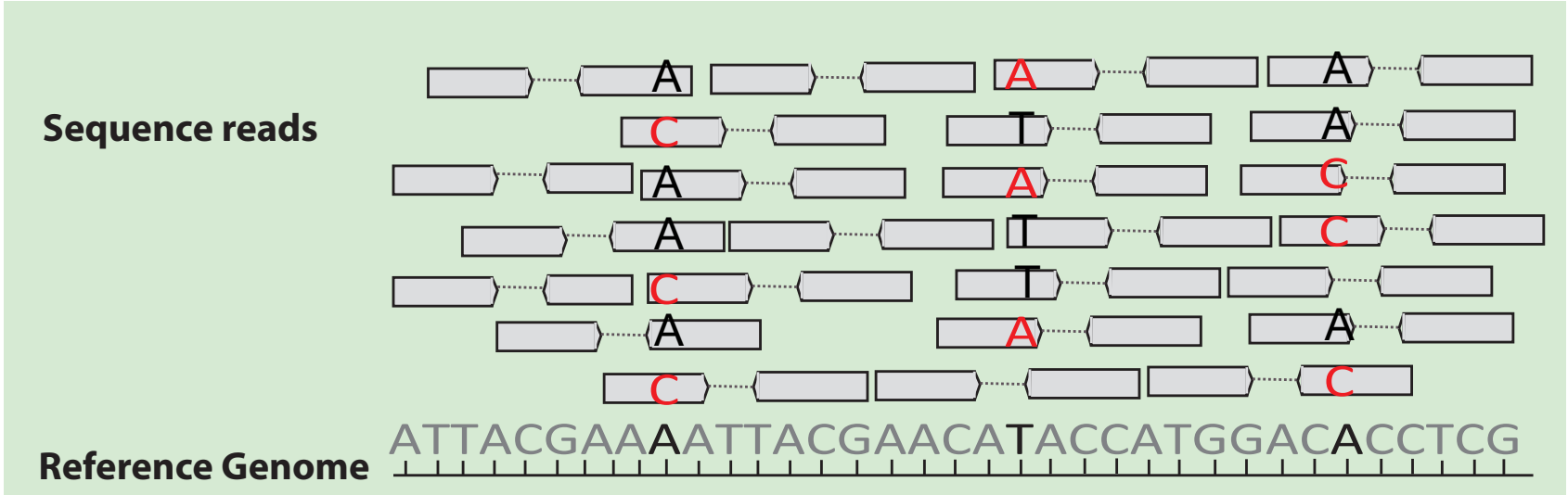
dbSNP (build 142): 112M SNPs, 88M "validated"

## 2. Somatic

- Acquired during cancer progression
- Not present in normal cells



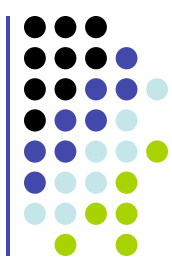
# Phasing



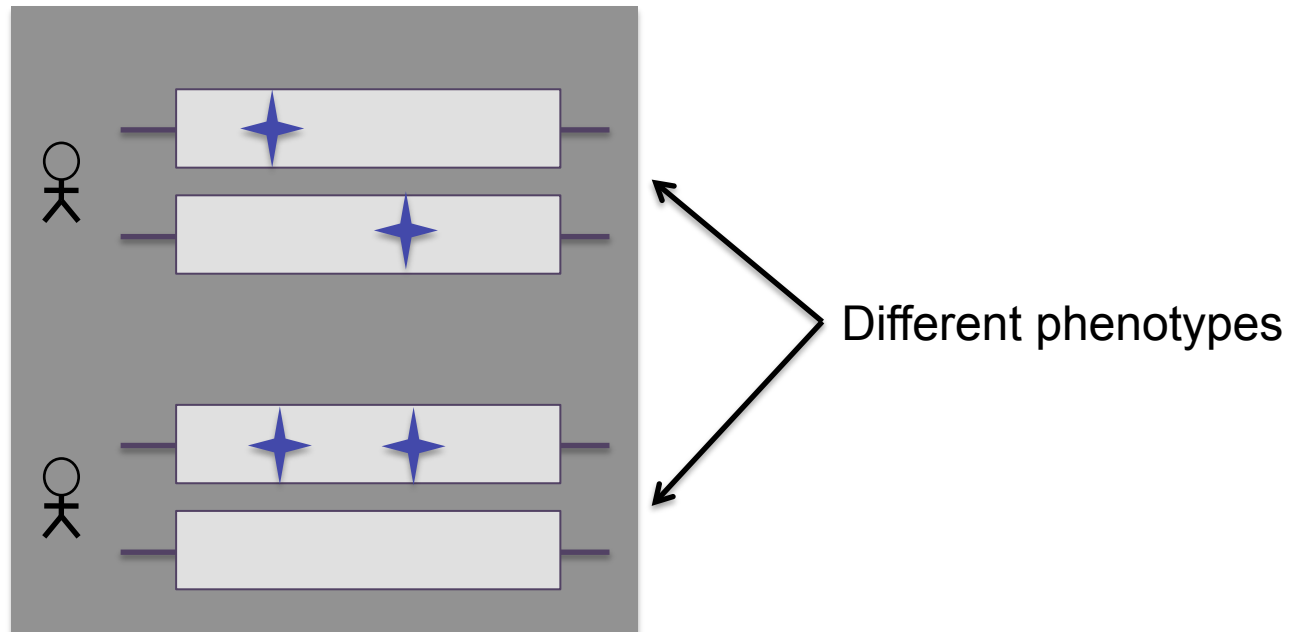
<b>Genotype</b>	A/C	A/T	A/C
-----------------	-----	-----	-----

<b>Haplotypes</b>	A	A	C
	C	T	A

**Phasing** is the process of recovering haplotypes from genotype data

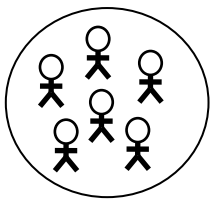


# Phasing – Compound Heterozygosity





# Phasing – Different Approaches



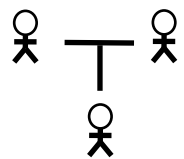
Population



Unrelated duos,  
trios



Population  
inference



Genetic



Pedigrees  
Trios



IBD analysis,  
inheritance state  
analysis



Molecular



Each physical read is a  
single molecule that can  
be directly phased



Read Assembly

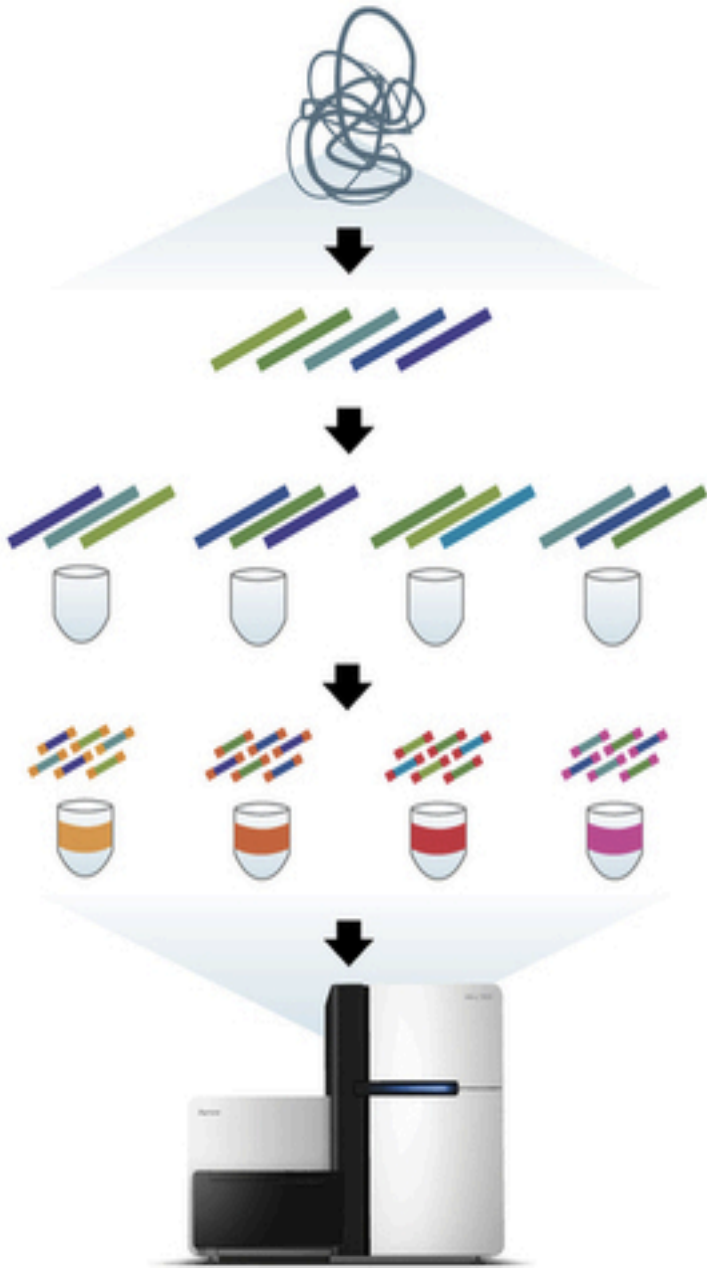
# Moleculo Overview

1. The sample DNA is sheared into fragments of about 10 kbp

2. The fragments are diluted and placed into 384 wells

3. Fragments are amplified through long-range PCR, cut into short fragments and barcoded

4. Short fragments are pooled together and sequenced





# Read Clouds



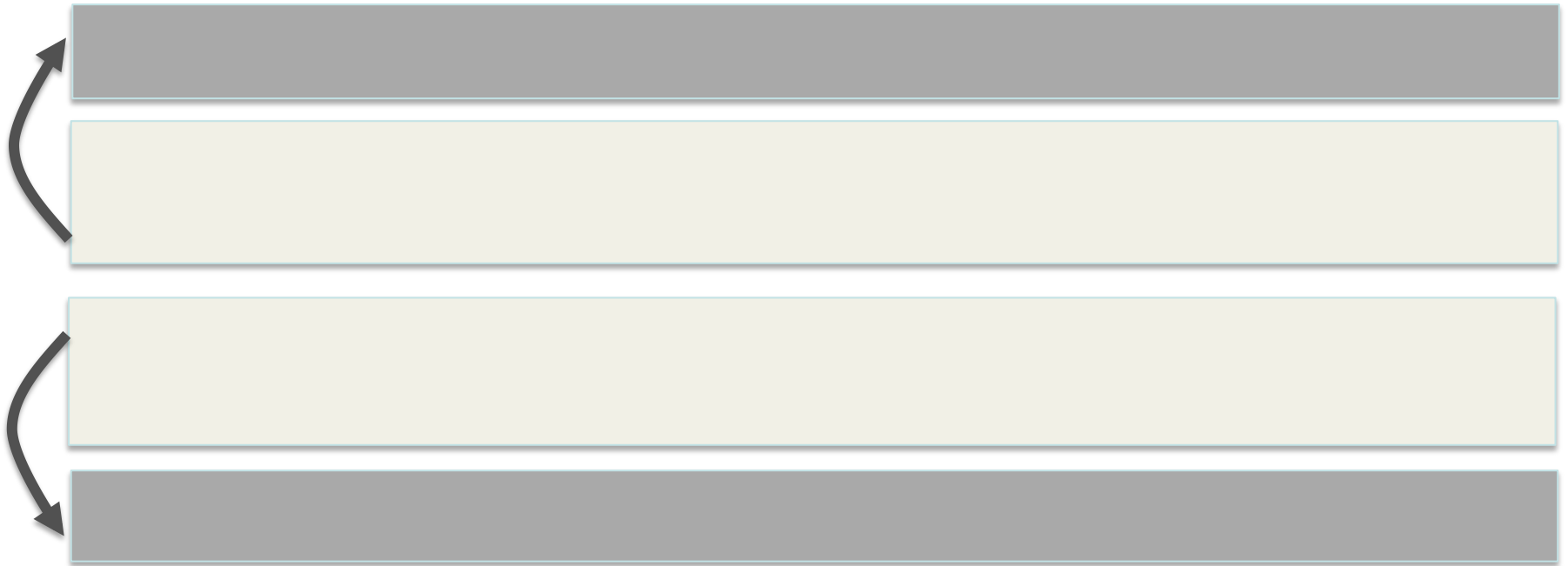
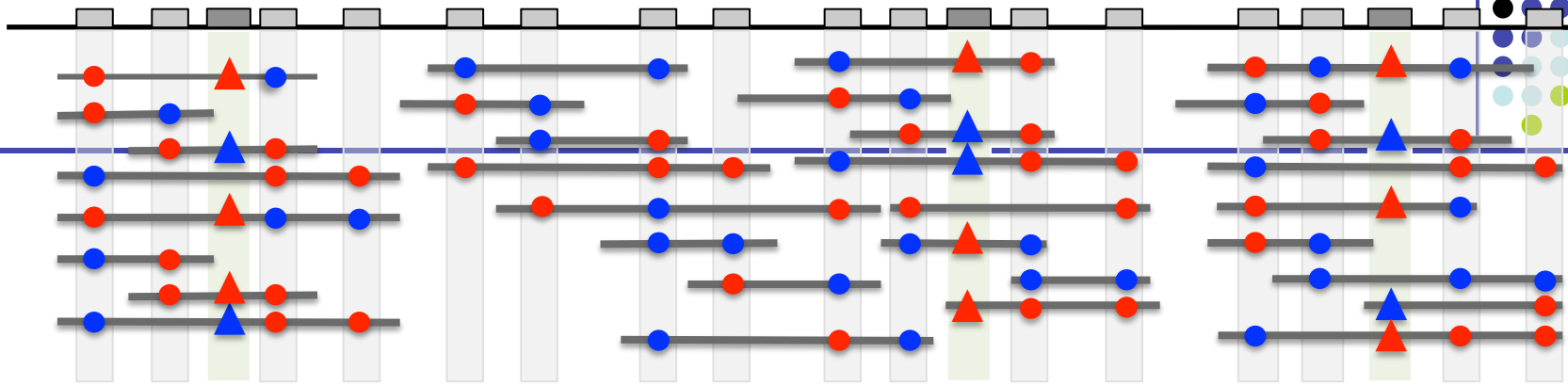
## Reference Genome

AACAGTAACCTTGATTACGTAAGTACCCCTTGACTAAAACCTCAAGGTACTGGATACCTGTAAACCRTCGAACTGAAACTAAAAGTAACTAACTAACTAACTAAGTAACTGACTAACTGTAACTGAAATGCAAACCTCCCG



A **read cloud** is a group of reads originated from one 10Kb fragment.

# Reference Genome

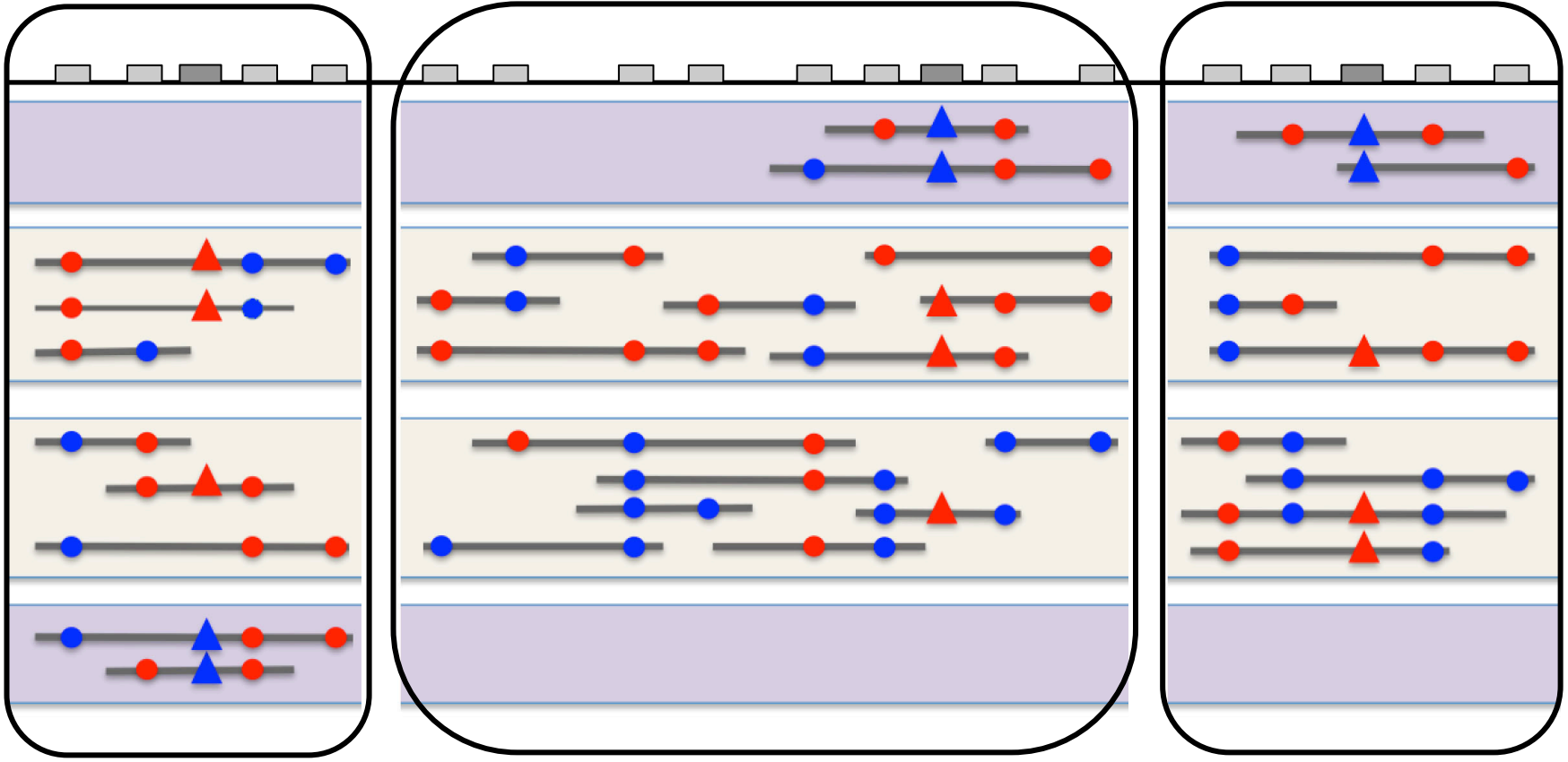


● Alternate Allele   ● Reference Allele   △ Somatic   ○ Germline



There can be **gaps** between some neighboring SNVs.

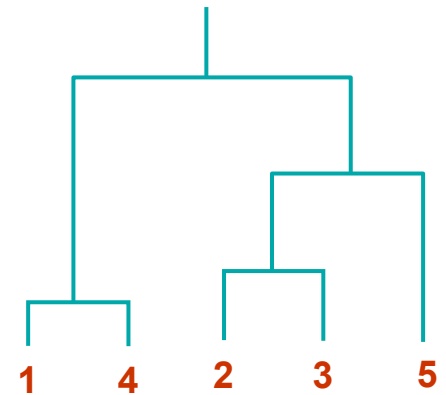
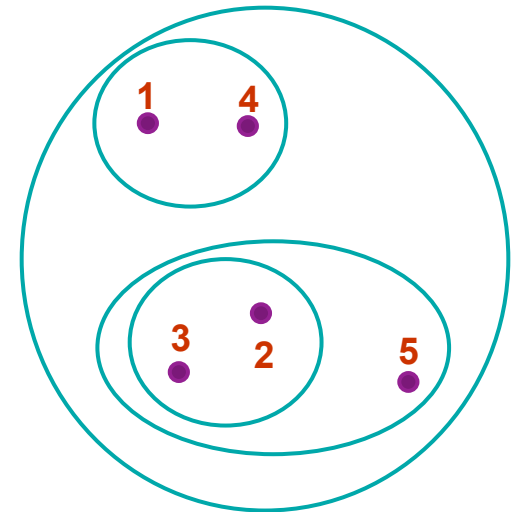
### Local Haplotype Blocks



● Alternate Allele   ● Reference Allele   △ Somatic   ○ Germline

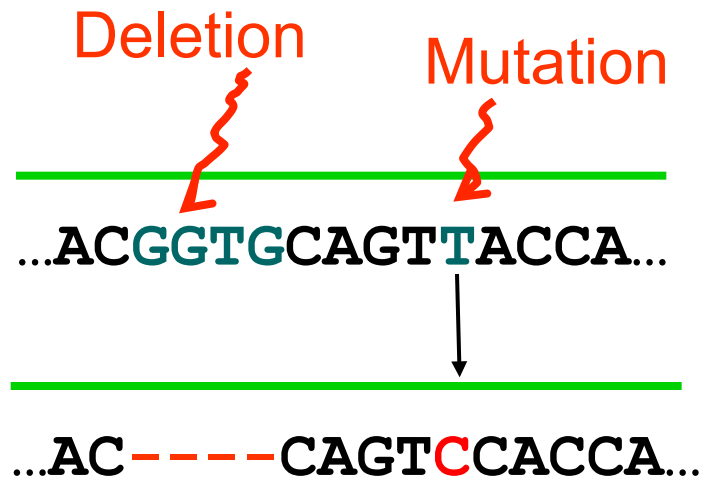


# Molecular Evolution and Phylogenetic Tree Reconstruction



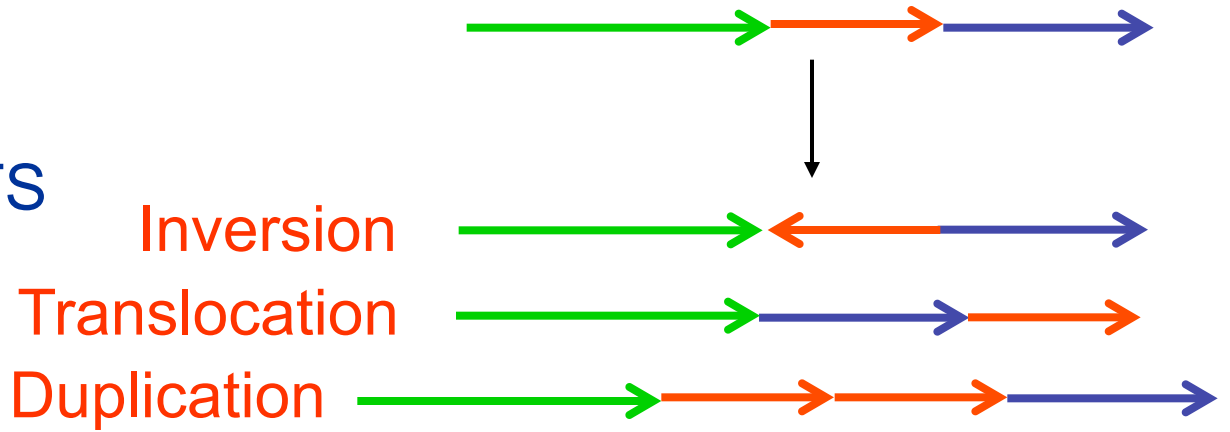


# Evolution at the DNA level



## SEQUENCE EDITS

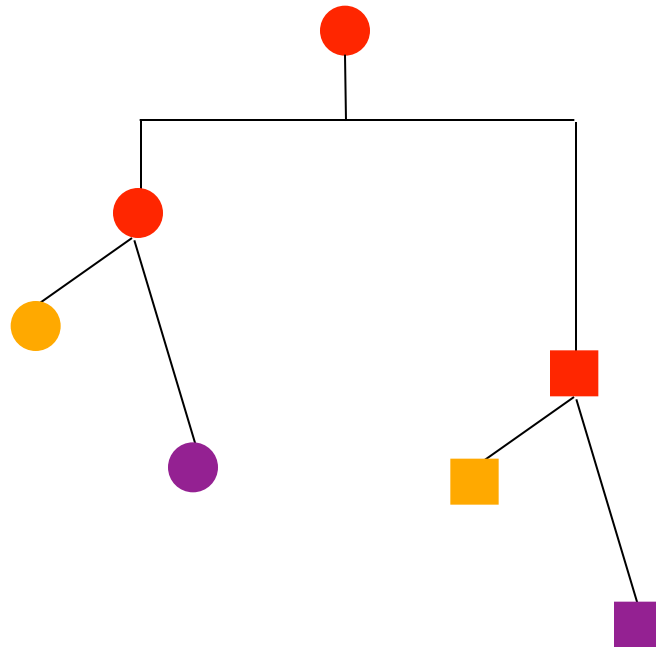
## REARRANGEMENTS





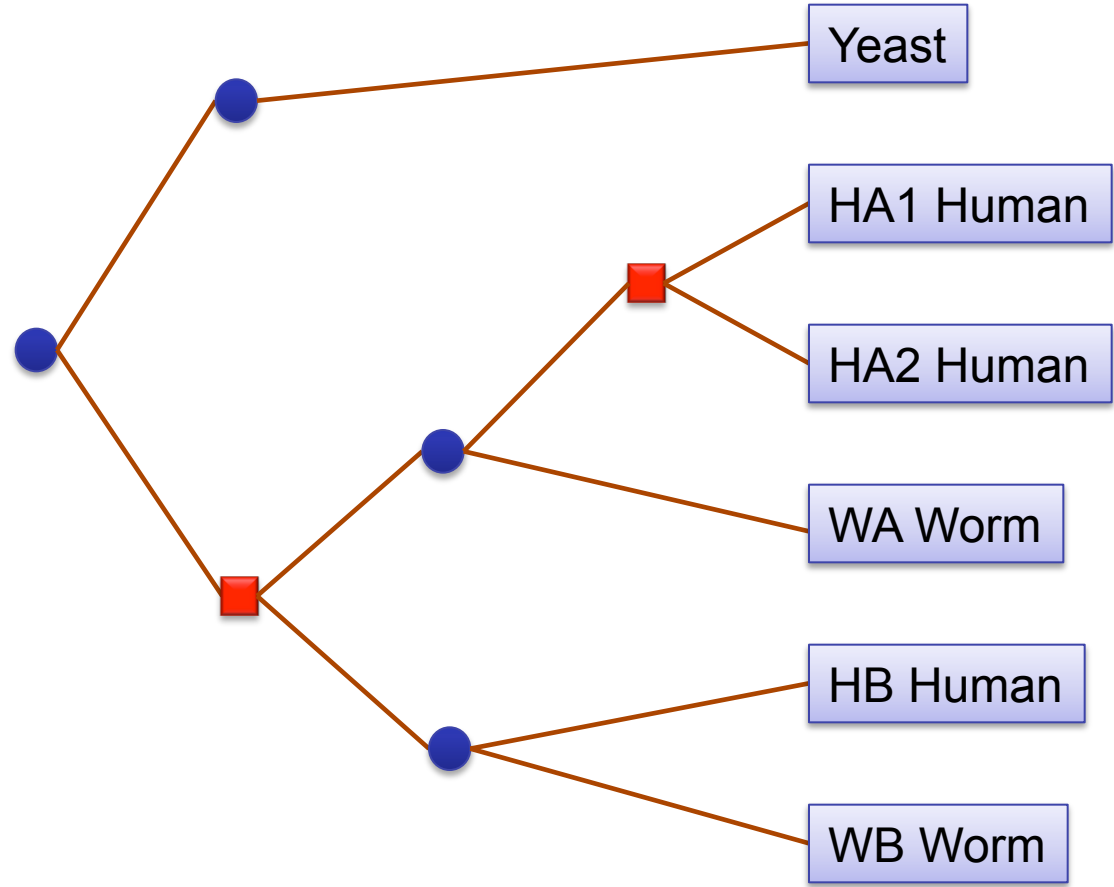
# Protein Phylogenies

- Proteins (genes) evolve by both duplication and species divergence





# Orthology and Paralogy



**Orthologs:**  
*Derived by  
speciation*

**Paralogs:**  
*Everything  
else*



# Orthology, Paralogy, Inparalogs, Outparalogs

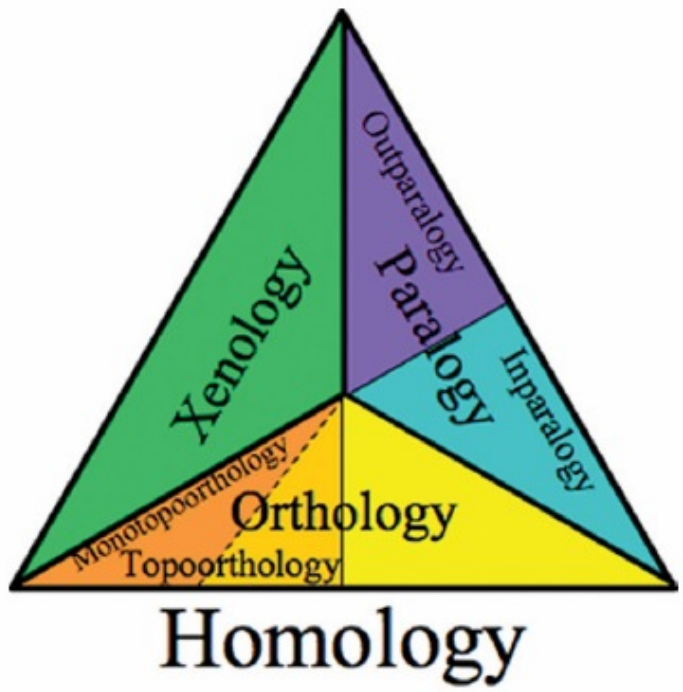


Figure 1. Refinements of homology.

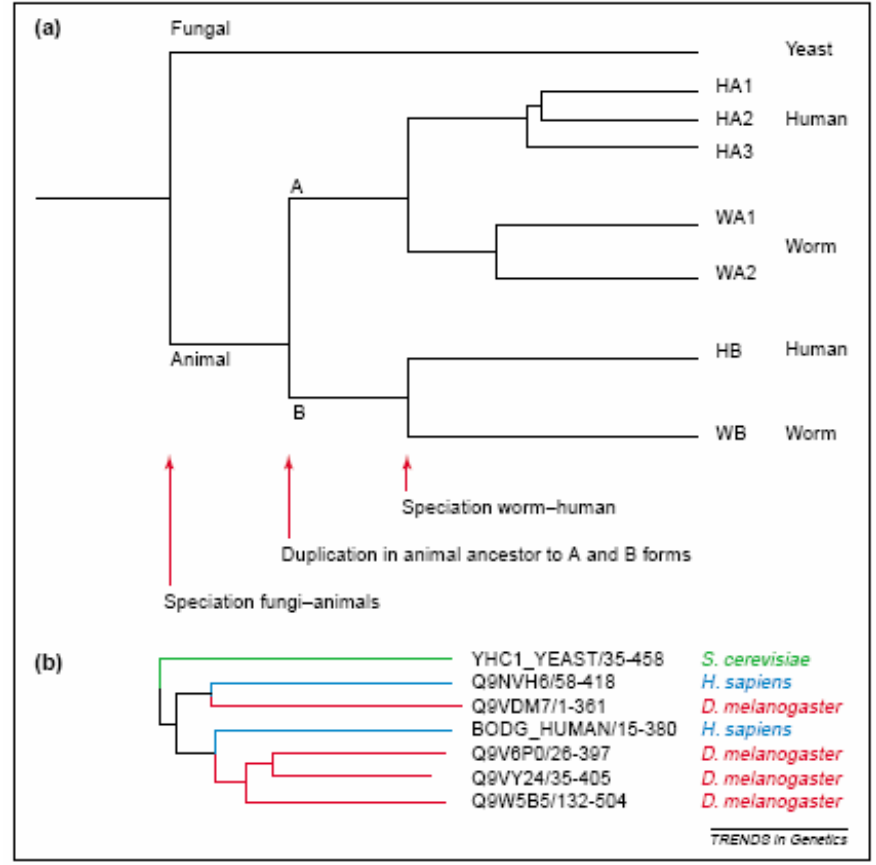


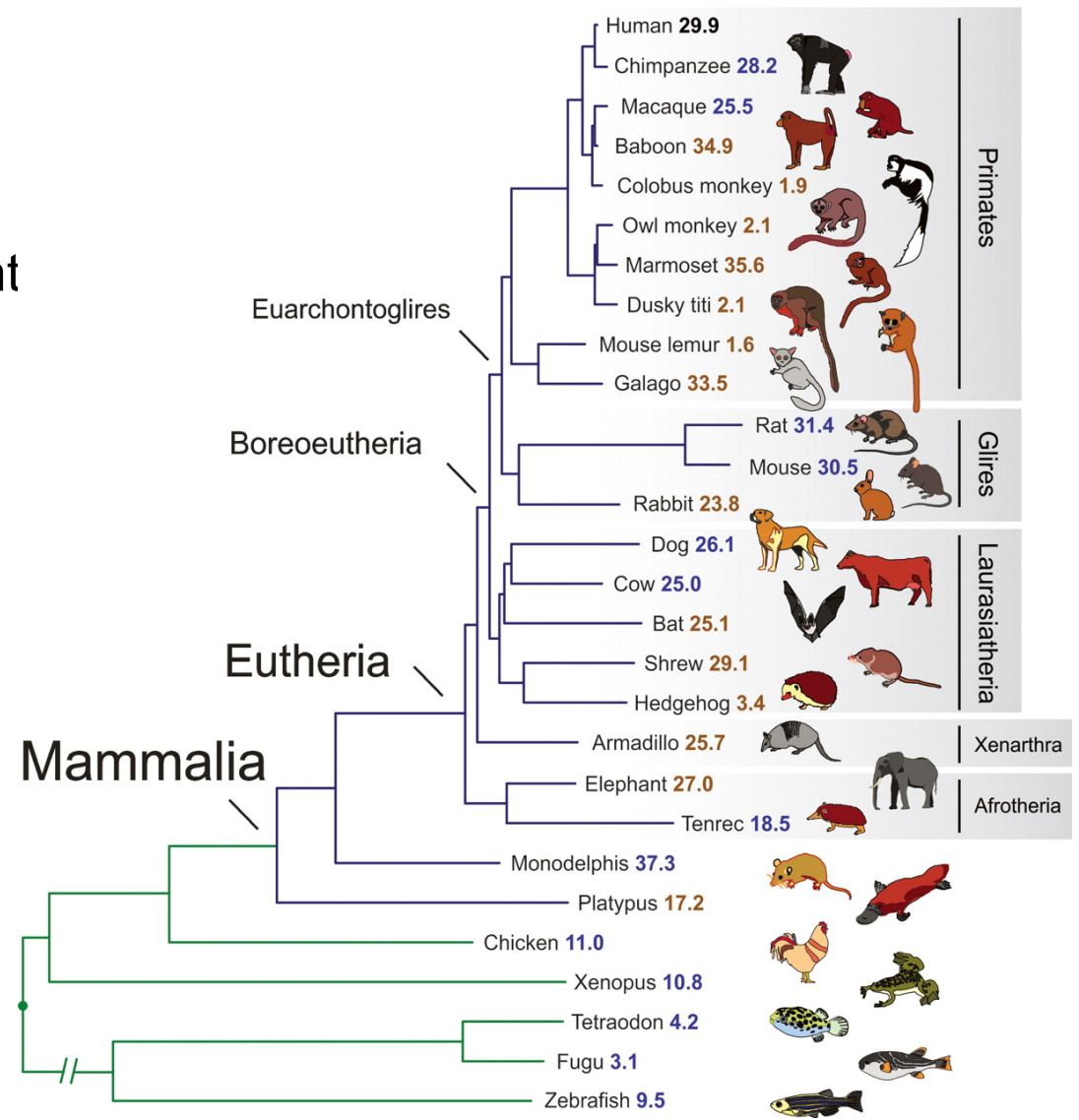
Fig. 1. The definition of inparalogs and outparalogs. (a) Consider an ancient gene inherited in the yeast, worm and human lineages. The gene was duplicated early in the animal lineage, before the human-worm split, into genes A and B. After the human-worm split, the A form was in turn duplicated independently in the human and worm lineages. In this scenario, the yeast gene is orthologous to all worm and human genes, which are all co-orthologous to the yeast gene. When comparing the human and worm genes, all genes in the HA\* set are co-orthologous to all genes in the WA\* set. The genes HA\* are hence 'inparalogs' to each other when comparing human to worm. By contrast, the genes HB and HA\* are 'outparalogs' when comparing human with worm. However, HB and HA\*, and WB and WA\* are inparalogs when comparing with yeast, because the animal-yeast split pre-dates the HA\*-HB duplication. (b) Real-life example of inparalogs:  $\gamma$ -butyrobetaine hydroxylases. The points of speciation and duplication are easily identifiable. The alignment is a subset of Pfam:PF03322 and the tree was generated by neighbor-joining in Belvu. All nodes have a bootstrap support exceeding 95%.





# Phylogenetic Trees

- Nodes: species
- Edges: time of independent evolution
- Edge length represents evolution time
  - AKA genetic distance
  - Not necessarily chronological time



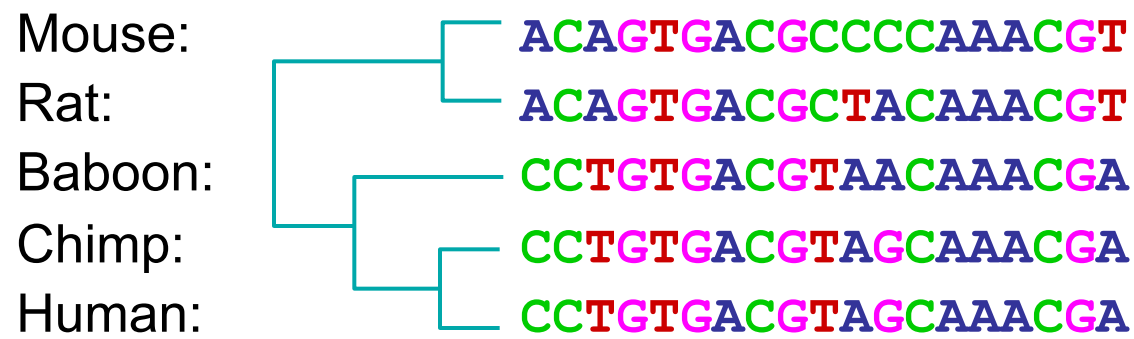


# Inferring Phylogenetic Trees

Trees can be inferred by several criteria:

- Morphology of the organisms
  - *Can lead to mistakes*
- Sequence comparison

**Example:**





# Inferring Phylogenetic Trees

- Sequence-based methods
  - Deterministic (Parsimony)
  - Probabilistic (SEMPHY)
- Distance-based methods
  - UPGMA
  - Neighbor-Joining
- Can compute distances from sequences



# Distance Between Two Sequences

## Basic principle:

- Distance proportional to degree of independent sequence evolution

Given sequences  $x^i$ ,  $x^j$ ,

$d_{ij}$  = distance between the two sequences

One possible definition:

$d_{ij}$  = fraction  $f$  of sites  $u$  where  $x^i[u] \neq x^j[u]$

Better scores are derived by modeling evolution as a continuous change process



# Molecular Evolution

## Modeling sequence substitution:

Consider what happens at a position for time  $\Delta t$ ,

- $P(t)$  = vector of probabilities of {A,C,G,T} at time  $t$
- $\mu_{AC}$  = rate of transition from A to C per unit time
- $\mu_A = \mu_{AC} + \mu_{AG} + \mu_{AT}$  rate of transition out of A
- $p_A(t+\Delta t) = p_A(t) - p_A(t) \mu_A \Delta t + p_C(t) \mu_{CA} \Delta t + p_G(t) \mu_{GA} \Delta t + p_T(t) \mu_{TA} \Delta t$



# Molecular Evolution

In matrix/vector notation, we get

$$P(t+\Delta t) = P(t) + Q P(t) \Delta t$$

where  $Q$  is the substitution rate matrix

$$Q = \begin{pmatrix} -\mu_A & \mu_{GA} & \mu_{CA} & \mu_{TA} \\ \mu_{AG} & -\mu_G & \mu_{CG} & \mu_{TG} \\ \mu_{AC} & \mu_{GC} & -\mu_C & \mu_{TC} \\ \mu_{AT} & \mu_{GT} & \mu_{CT} & -\mu_T \end{pmatrix}$$



# Molecular Evolution

- This is a differential equation:

$$P'(t) = Q P(t)$$

- $Q \Rightarrow$  prob. distribution over  $\{A, C, G, T\}$  at each position, stationary (equilibrium) frequencies  $\pi_A, \pi_C, \pi_G, \pi_T$
- Each  $Q$  is an evolutionary model
  - Some work better than others



# Evolutionary Models

- Jukes-Cantor

$$Q = \begin{pmatrix} * & \frac{\mu}{4} & \frac{\mu}{4} & \frac{\mu}{4} \\ \frac{\mu}{4} & * & \frac{\mu}{4} & \frac{\mu}{4} \\ \frac{\mu}{4} & \frac{\mu}{4} & * & \frac{\mu}{4} \\ \frac{\mu}{4} & \frac{\mu}{4} & \frac{\mu}{4} & * \end{pmatrix}$$

- Kimura

$$Q = \begin{pmatrix} * & \kappa & 1 & 1 \\ \kappa & * & 1 & 1 \\ 1 & 1 & * & \kappa \\ 1 & 1 & \kappa & * \end{pmatrix}$$

- Felsenstein

$$Q = \begin{pmatrix} * & \pi_T & \pi_T & \pi_T \\ \pi_C & * & \pi_C & \pi_C \\ \pi_A & \pi_A & * & \pi_A \\ \pi_G & \pi_G & \pi_G & * \end{pmatrix}$$

- HKY

$$Q = \begin{pmatrix} * & \kappa\pi_T & \pi_T & \pi_T \\ \kappa\pi_C & * & \pi_C & \pi_C \\ \pi_A & \pi_A & * & \kappa\pi_A \\ \pi_G & \pi_G & \kappa\pi_G & * \end{pmatrix}$$





# Estimating Distances

- Solve the differential equation and compute expected evolutionary time given sequences

$$P'(t) = Q P(t)$$

Jukes-Cantor:

$$\text{Let } P_{AA}(t) = P_{CC}(t) = P_{GG}(t) = P_{TT}(t) = r$$

$$P_{AC}(t) = \dots = P_{TG}(t) = s$$

Then,

$$r'(t) = -\frac{3}{4} r(t) \mu + \frac{3}{4} s(t) \mu$$

$$s'(t) = -\frac{1}{4} s(t) \mu + \frac{1}{4} r(t) \mu$$

Which is satisfied by

$$r(t) = \frac{1}{4} (1 + 3e^{-\mu t})$$

$$s(t) = \frac{1}{4} (1 - e^{-\mu t})$$



# Estimating Distances

- Solve the differential equation and compute expected evolutionary time given sequences

$$P'(t) = Q P(t)$$

Jukes-Cantor:

$$P = \begin{pmatrix} \frac{1}{4} + \frac{3}{4}e^{-t\mu} & \frac{1}{4} - \frac{1}{4}e^{-t\mu} & \frac{1}{4} - \frac{1}{4}e^{-t\mu} & \frac{1}{4} - \frac{1}{4}e^{-t\mu} \\ \frac{1}{4} - \frac{1}{4}e^{-t\mu} & \frac{1}{4} + \frac{3}{4}e^{-t\mu} & \frac{1}{4} - \frac{1}{4}e^{-t\mu} & \frac{1}{4} - \frac{1}{4}e^{-t\mu} \\ \frac{1}{4} - \frac{1}{4}e^{-t\mu} & \frac{1}{4} - \frac{1}{4}e^{-t\mu} & \frac{1}{4} + \frac{3}{4}e^{-t\mu} & \frac{1}{4} - \frac{1}{4}e^{-t\mu} \\ \frac{1}{4} - \frac{1}{4}e^{-t\mu} & \frac{1}{4} - \frac{1}{4}e^{-t\mu} & \frac{1}{4} - \frac{1}{4}e^{-t\mu} & \frac{1}{4} + \frac{3}{4}e^{-t\mu} \end{pmatrix}$$



# Estimating Distances

Let  $p$  = probability a base is different between two sequences,  
Solve to find  $t$

- Jukes-Cantor  $r(t) = 1 - p = \frac{1}{4} (1 + 3e^{-\mu t})$

$$p = \frac{3}{4} - \frac{3}{4} e^{-\mu t}$$

$$\frac{3}{4} - p = \frac{3}{4} e^{-\mu t}$$

$$1 - 4p/3 = e^{-\mu t}$$

Therefore,

$$\mu t = -\ln(1 - 4p/3)$$

Letting

$d = \frac{3}{4} \mu t$ , denoting substitutions per site,

$$d = -\frac{3}{4} \ln\left(1 - \frac{4}{3}p\right)$$



# Estimating Distances

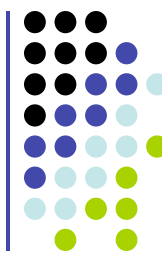
$d$ : Branch length in terms of substitutions per site

- Jukes-Cantor

$$d = -\frac{3}{4} \ln\left(1 - \frac{4}{3}p\right)$$

- Kimura

$$d = -\frac{1}{2} \ln(1 - 2P - Q) - \frac{1}{4} \ln(1 - 2Q)$$



# Simple method for building tree: UPGMA

UPGMA (unweighted pair group method using arithmetic averages)  
Or the **Average Linkage Method**

Given two disjoint clusters  $C_i$ ,  $C_j$  of sequences,

$$d_{ij} = \frac{1}{|C_i| \times |C_j|} \sum_{\{p \in C_i, q \in C_j\}} d_{pq}$$

Claim that if  $C_k = C_i \cup C_j$ , then distance to another cluster  $C_l$  is:

$$d_{kl} = \frac{d_{il} |C_i| + d_{jl} |C_j|}{|C_i| + |C_j|}$$



# Algorithm: Average Linkage

## Initialization:

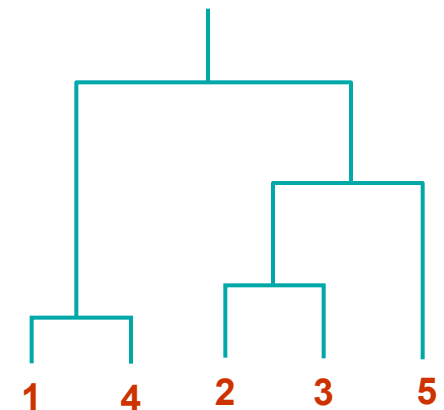
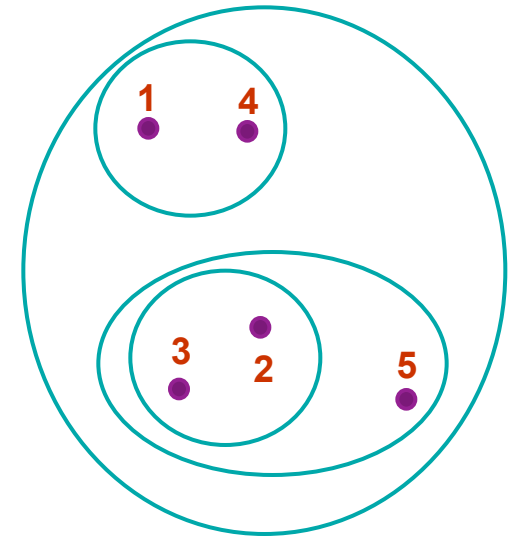
Assign each  $x_i$  into its own cluster  $C_i$   
Define one leaf per sequence, height 0

## Iteration:

Find two clusters  $C_i, C_j$  s.t.  $d_{ij}$  is min  
Let  $C_k = C_i \cup C_j$   
Define node connecting  $C_i, C_j$ , and place it at  
height  $d_{ij}/2$   
Delete  $C_i, C_j$

## Termination:

When two clusters  $i, j$  remain, place root at  
height  $d_{ij}/2$





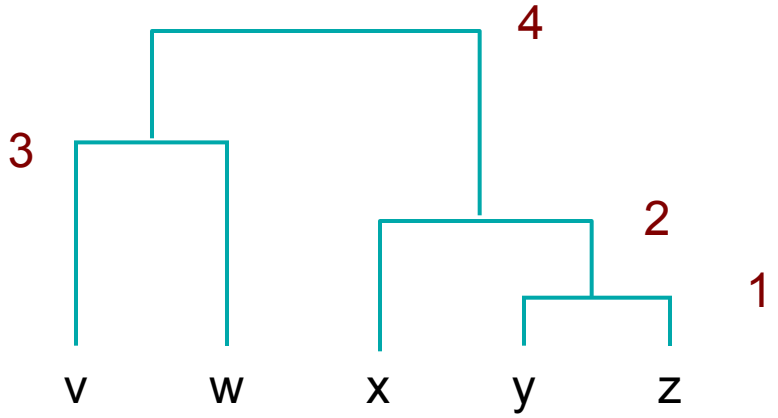
# Average Linkage Example

	v	w	x	y	z
v	0	6	8	8	8
w		0	8	8	8
x			0	4	4
y				0	2
z					0

	v	w	xyz
v	0	6	8
w		0	8
xyz			0

	vw	xyz
vw	0	8
xyz		0

	v	w	x	yz
v	0	6	8	8
w		0	8	8
x			0	4
yz				0



# Ultrametric Distances and Molecular Clock



## Definition:

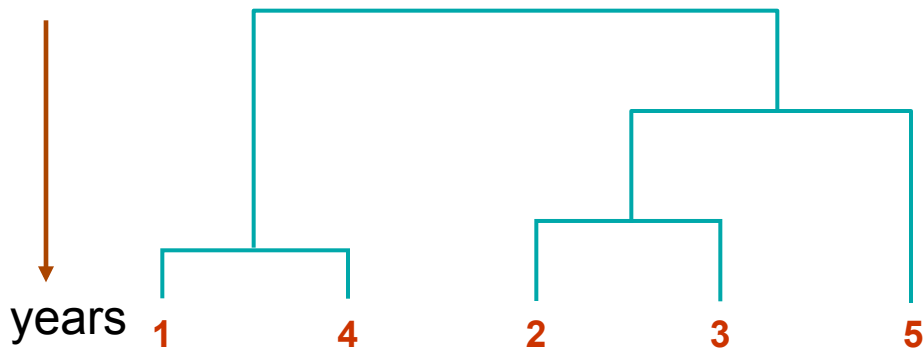
A distance function  $d(.,.)$  is ultrametric if for any three distances  $d_{ij} \leq d_{ik} \leq d_{ij}$ , it is true that

$$d_{ij} \leq d_{ik} = d_{jk}$$

## The Molecular Clock:

The evolutionary distance between species  $x$  and  $y$  is  $2 \times$  the Earth time to reach the nearest common ancestor

That is, the molecular clock has constant rate in all species

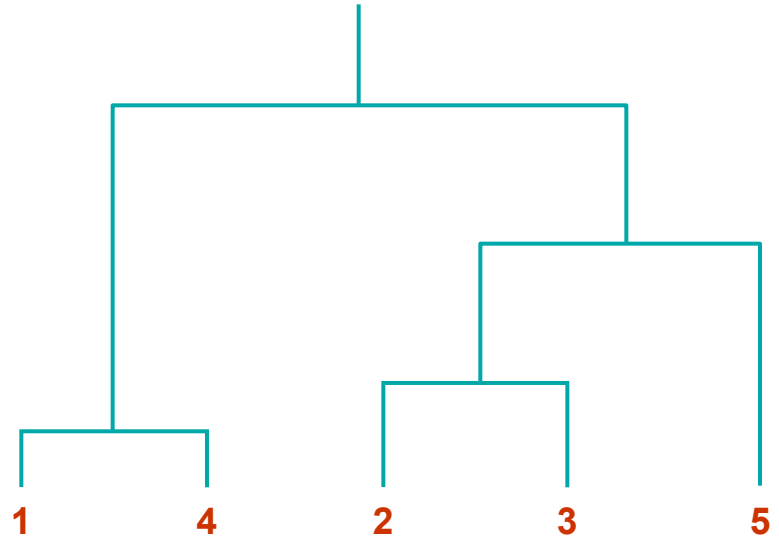


The molecular clock results in ultrametric distances





# Ultrametric Distances & Average Linkage



Average Linkage is guaranteed to reconstruct correctly a binary tree with ultrametric distances

Proof: Exercise

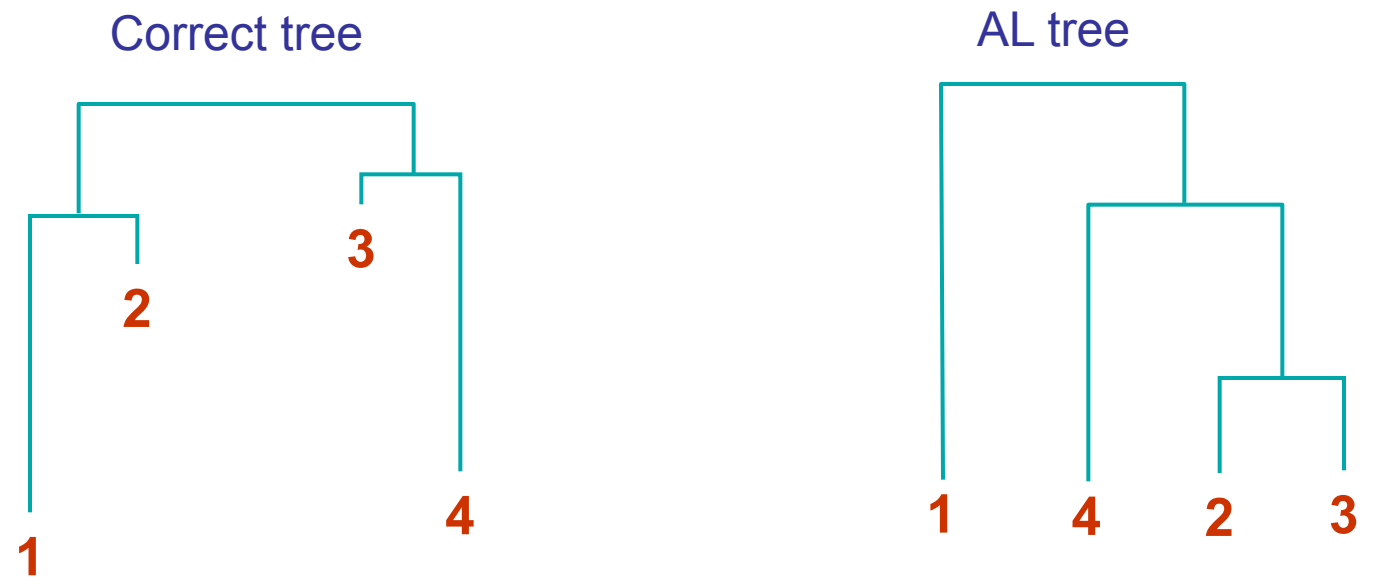


# Weakness of Average Linkage

Molecular clock: all species evolve at the same rate (Earth time)

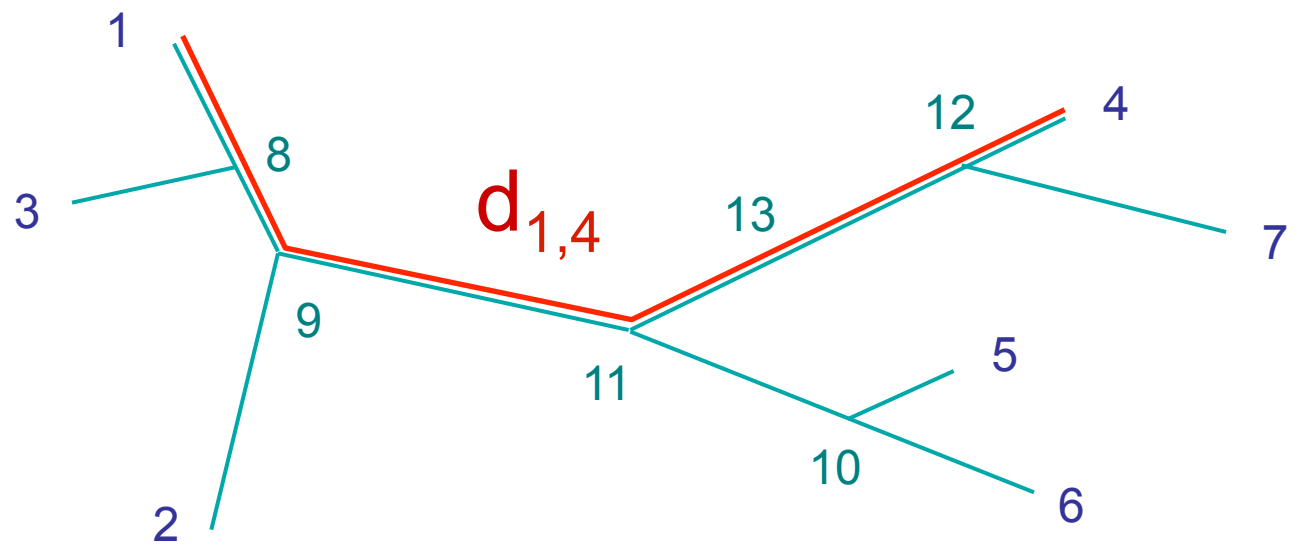
However, certain species (e.g., mouse, rat) evolve much faster

Example where UPGMA messes up:





# Additive Distances



Given a tree, a distance measure is **additive** if the distance between any pair of leaves is the sum of lengths of edges connecting them

Given a tree  $T$  & additive distances  $d_{ij}$ , can uniquely reconstruct edge lengths:

- Find two neighboring leaves  $i, j$ , with common parent  $k$
- Place parent node  $k$  at distance  $d_{km} = \frac{1}{2} (d_{im} + d_{jm} - d_{ij})$  from any node  $m \neq i, j$



# Additive Distances



For any four leaves  $x, y, z, w$ , consider the three sums

$$\begin{aligned} & d(x, y) + d(z, w) \\ & d(x, z) + d(y, w) \\ & d(x, w) + d(y, z) \end{aligned}$$

One of them is smaller than the other two, which are equal

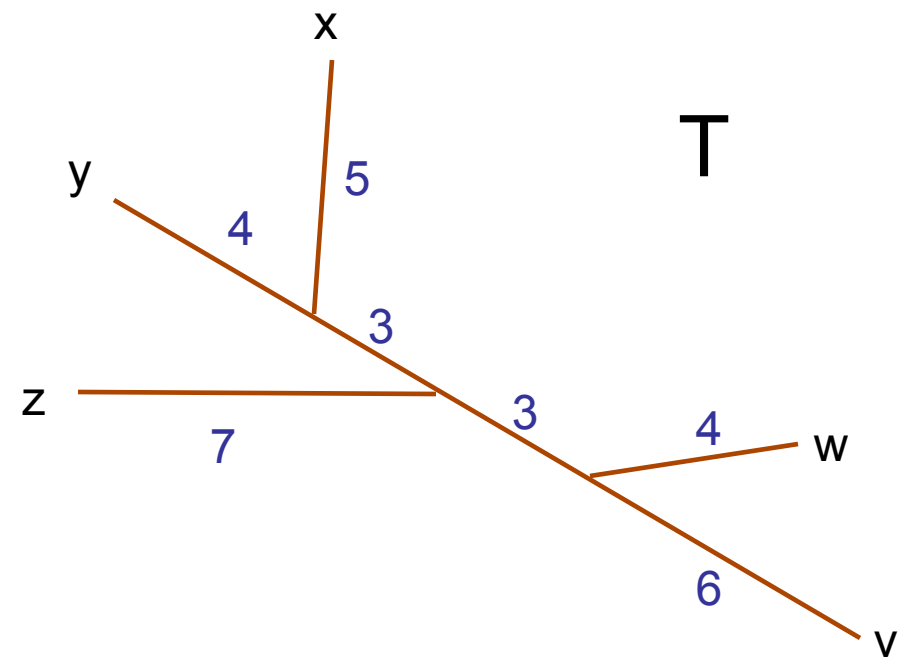
$$d(x, y) + d(z, w) < d(x, z) + d(y, w) = d(x, w) + d(y, z)$$



# Reconstructing Additive Distances Given T

D

	v	w	x	y	z
v	0	10	17	16	16
w		0	15	14	14
x			0	9	15
y				0	14
z					0



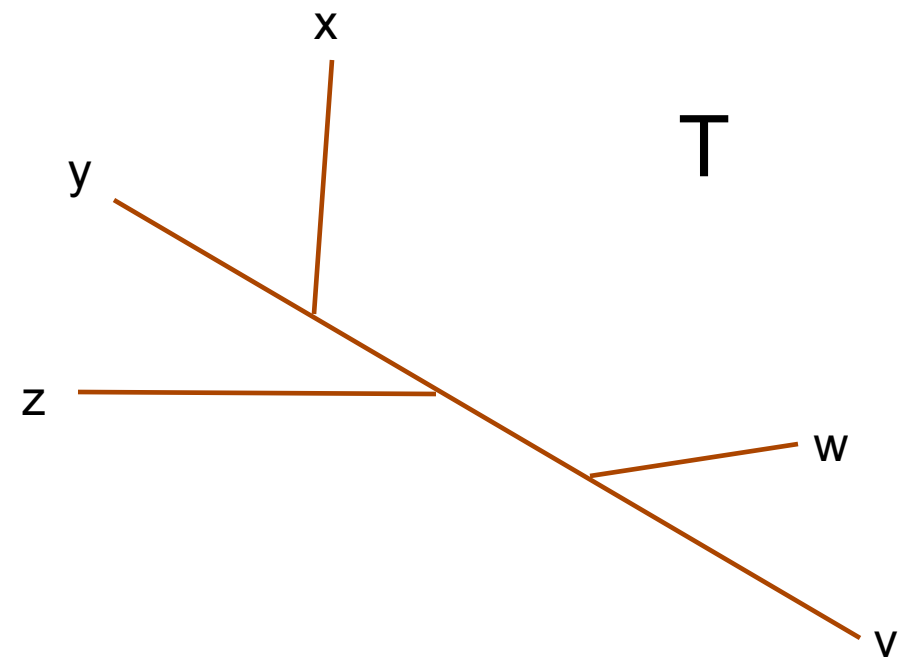
If we know T and D, but do not know the length of each leaf, we can reconstruct those lengths



# Reconstructing Additive Distances Given T

D

	v	w	x	y	z
v	0	10	17	16	16
w		0	15	14	14
x			0	9	15
y				0	14
z					0

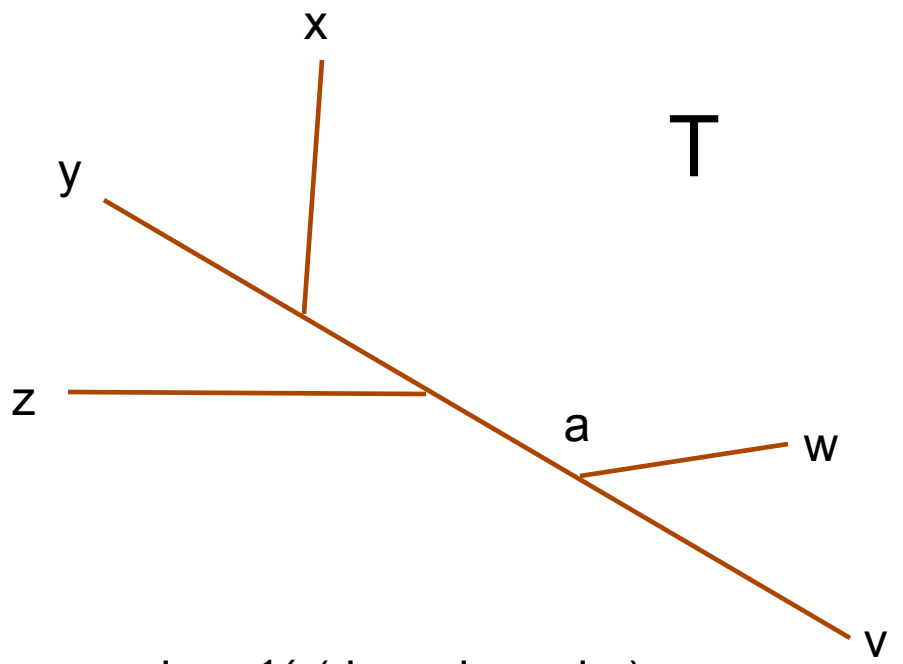




# Reconstructing Additive Distances Given T

D

	v	w	x	y	z
v	0	10	17	16	16
w		0	15	14	14
x			0	9	15
y				0	14
z					0



D<sub>1</sub>

	a	x	y	z
a	0	11	10	10
x		0	9	15
y			0	14
z				0

$$d_{ax} = \frac{1}{2} (d_{vx} + d_{wx} - d_{vw})$$

$$d_{ay} = \frac{1}{2} (d_{vy} + d_{wy} - d_{vw})$$

$$d_{az} = \frac{1}{2} (d_{vz} + d_{wz} - d_{vw})$$



# Reconstructing Additive Distances Given T

$D_1$

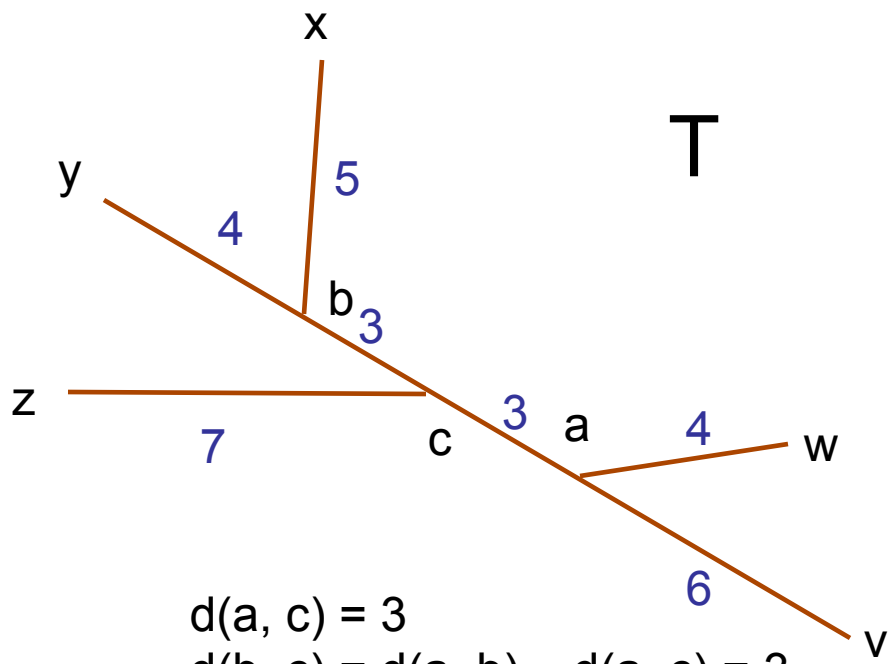
	a	x	y	z
a	0	11	10	10
x		0	9	15
y			0	14
z				0

$D_2$

	a	b	z
a	0	6	10
b		0	10
z			0

$D_3$

	a	c
a	0	3
c		0



- $d(a, c) = 3$
- $d(b, c) = d(a, b) - d(a, c) = 3$
- $d(c, z) = d(a, z) - d(a, c) = 7$
- $d(b, x) = d(a, x) - d(a, b) = 5$
- $d(b, y) = d(a, y) - d(a, b) = 4$
- $d(a, w) = d(z, w) - d(a, z) = 4$
- $d(a, v) = d(z, v) - d(a, z) = 6$

**Correct!!!**





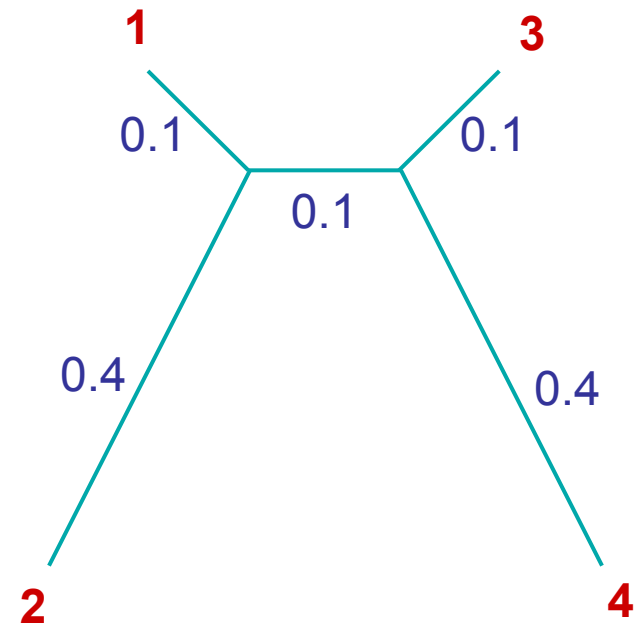
# Neighbor-Joining

- Guaranteed to produce the correct tree if distance is additive
- May produce a good tree even when distance is not additive

## Step 1: Finding neighboring leaves

Define

$$D_{ij} = (N - 2) d_{ij} - \sum_{k \neq i} d_{ik} - \sum_{k \neq j} d_{jk}$$

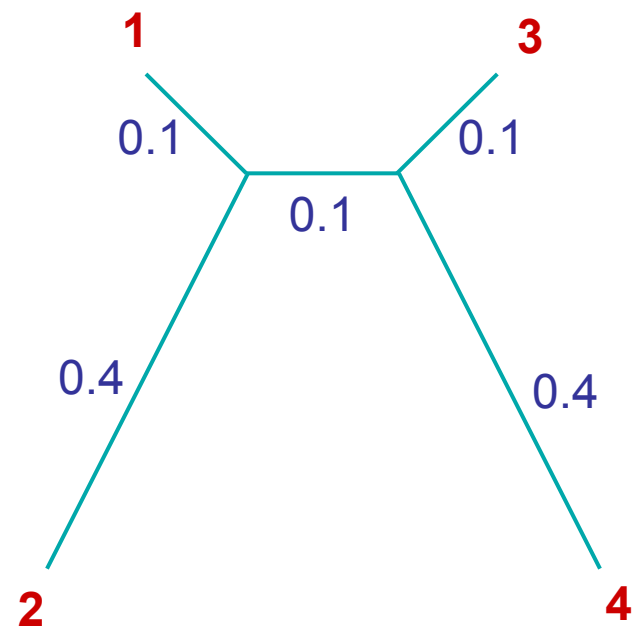
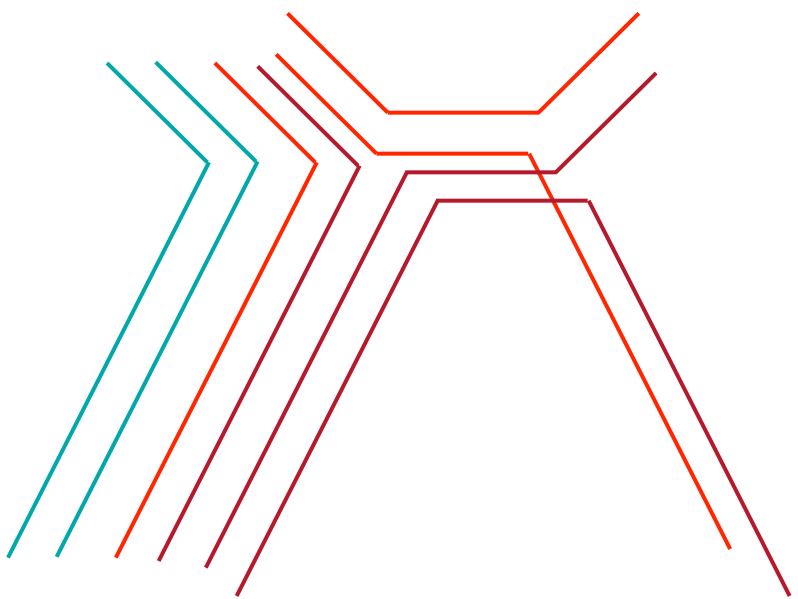


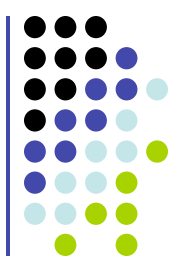
**Claim:** The above “magic trick” ensures that  $i, j$  are neighbors if  $D_{ij}$  is minimal



# Neighbor-Joining

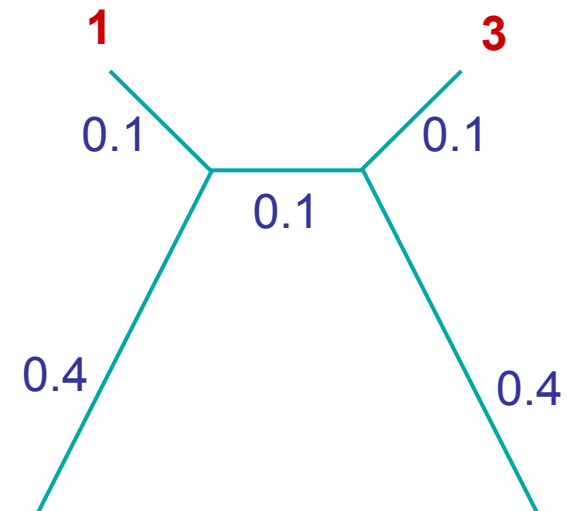
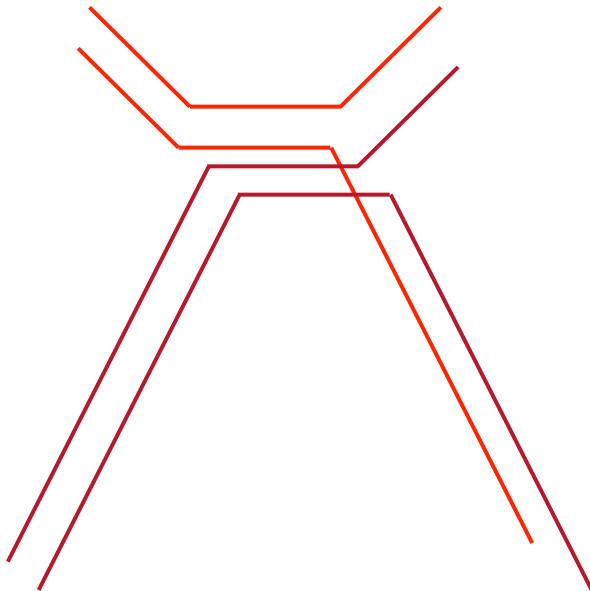
$$D_{ij} = (N - 2) d_{ij} - \sum_{k \neq i} d_{ik} - \sum_{k \neq j} d_{jk}$$





# Neighbor-Joining

$$D_{ij} = (N - 2) d_{ij} - \sum_{k \neq i} d_{ik} - \sum_{k \neq j} d_{jk}$$



- All leaf edges appear negatively exactly twice
- All other edges appear negatively once for every path from each of the two leaves  $i, j$ , to leaves  $k \neq i, j$