

Class 4: Agenda, Questions, and Links

1 Warm-Up

Suppose you flip a p -biased coin n times. What does Markov's inequality tell you about the probability that you see more than $2pn$ heads? What does Chebyshev's inequality tell you about that probability?

2 Questions?

Any questions from the minilectures or warmup? (Markov and Chebyshev's inequalities).

3 Sampling-Based Median

Sampling-Based Median Algorithm

Median:(A list S of n distinct numbers, where n is odd):

1. Let $t = n^{3/4}$. Sample $R = \{r_1, \dots, r_t\} \subseteq S$ by drawing r_i uniformly at random, independently.
2. Sort R in time $O(t \log t)$. Henceforth, assume that $r_1 \leq r_2 \leq \dots \leq r_t$.
3. Let $a = r_{t/2-\sqrt{n}}$, $b = r_{t/2+\sqrt{n}}$.
4. Let $N_{<a}$ and $N_{>b}$ denote the number of elements in S less than a and greater than b respectively.
5. Let $T = \{x \in S : a \leq x \leq b\}$. Construct T , and compute $N_{<a}$ and $N_{>b}$, in time $O(n)$.
6. If $|T| < 4t$, sort T in time $O(t \log t)$; otherwise output FAIL.
7. If $N_{<a}, N_{>b} \leq n/2$ (aka, $median(S) \in T$):
 - Return the i 'th smallest element of T , where $i = (n + 1)/2 - N_{<a}$.
8. Otherwise, output FAIL.

[We'll see an example on a slide; this slide is posted on the course website.]

[**Note:** Above there should be some floors or ceilings or something. Don't worry about it, and ignore off-by-one errors throughout this class.]

4 Analyzing the sampling-based median algorithm

You will analyze this algorithm in group work.

Group Work

Note: Throughout this group work, don't worry about \leq vs $<$, or whether or not something is true up to ± 1 , or anything small like that.

1. Make sure that you all understand the algorithm. Pseudo-code is above, and the one-slide example is available on the course website, in the class-by-class resources for Class 4. Ask/answer any questions that you have amongst yourselves, and flag down a member of the course staff if you still have questions.
2. Suppose that you could show that:
 - with probability ≥ 0.9 , the median of S is in the list T ; and
 - with probability ≥ 0.9 , $|T| < 4t$.

Explain (to each other) why these two things would imply that the algorithm returns the correct answer with probability ≥ 0.8 . And if it does not return the median then it returns FAIL.

3. In the following parts, you will show that the median of S is in T , with probability at least 0.9. Let m be the median of S . Consider two events:
 - A. $|\{r_i \in R : r_i < m\}| > \frac{t}{2} - \sqrt{n}$
 - B. $|\{r_i \in R : r_i > m\}| > \frac{t}{2} - \sqrt{n}$
 - (a) Explain why, if both of these events hold, then $\text{median}(S) \in T$.
 - (b) Use Chebyshev's inequality to bound the probability that the first event does not hold. (Hint: let X_i be the indicator random variable that is 1 iff $r_i < m$, and consider $\sum_i X_i$).
 - (c) Convince yourself that the same argument will work for the second event, and write a statement of the form:

$$\Pr[\text{median}(S) \in T] \geq 1 - \text{----}.$$

4. Now, we turn our attention to the probability that $|T| < 4t$.
 - (a) Explain why it is sufficient to show that a is *not* one of the smallest $n/2 - 2t$ elements of S , and b is *not* one of the largest $n/2 + 2t$ elements of S .
 - (b) Use Chebyshev's inequality to bound the probability that a is not one of the smallest $n/2 - 2t$ elements of S . (Hint: Consider the indicator random variable Y_i that is 1 if r_i is in the smallest $n/2 - 2t$ elements of S . Argue that a is one

of the smallest $n/2 - 2t$ elements of S iff $\sum_i Y_i \geq t/2 - \sqrt{n}$ (why?) and apply Chebyshev's inequality.)

- (c) Convince yourself that the analogous statement for b , and write a statement of the form:

$$\Pr[|T| < 4t] \geq 1 - \text{-----}.$$

5. Put it all together! At this point, you've done all the work to prove that the algorithm succeeds with probability at least $1 - \text{-----}$. Make sure you understand how the pieces fit together, and fill in the last blank.
6. **(Bonus, if time)** This algorithm runs in time $O(n)$. What is the leading constant in that big-Oh notation, assuming "sample a random element of S " and comparing two numbers are each a single operation? This constant depends on how you implement each step – what's the best you can do? Can you come up with a lower bound on the number of (expected or worst-case) comparisons needed to compute the median? How close is your answer above?
7. **(Bonus, if time)** We made a lot of decisions in the above algorithm/proof. For example, we chose $t = n^{3/4}$. We chose a and b to be $t/2 \pm \sqrt{n}$, but we could have chosen them to be $t/2 \pm \Delta$ for some parameter Δ (that may depend on n). We chose that the algorithm would output FAIL if $|T| < 4t$, but we could have said FAIL if $|T| < \alpha t$ for some parameter α (that could possibly also depend on n). Why did we make the decisions that we made? Are there other choices of (t, Δ, α) that would work? Are there some that would work better? Play around with these and see what the requirements are!