## **PEAKSTREAM**

# The PeakStream Platform for Many-Core Computing

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Engineering Director
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## **PeakStream History**

#### » PeakStream

- Startup company
- Founded February 2005
- 35 people
- Based in silicon valley

#### » PeakStream Mission Statement

 Provide a software platform for High Performance Computing that unlocks the power of a new generation of processors, from GPUs to multi-core CPUs

#### The PeakStream Team

#### » Founder: Matthew Papakipos

- Former NVIDIA Director of GPU Architecture: NV20 & NV40 Lead, XBox
- Graphics software standards: OpenGL & DirectX
- Supercomputers: MasPar & Connection Machine

#### » Chief Scientist: Pat Hanrahan

- Stanford computer science professor
- Led the Brook project (more on this later)

#### » Brian Grant

- Software architect, compiler expert
- Formerly at Transmeta

#### » Chris Demetriou

- Software architect, systems expert
- Formerly at SiByte/Broadcom, NetBSD

## Google & PeakStream

- » PeakStream was acquired by Google in May, 2007
  - Existing product line sales were discontinued
  - PeakStream's future is as part of Google
- » This presentation is a bit of history
  - The founding of the PeakStream
  - The technology
  - The product
  - The Stanford connection

## Before PeakStream: Setting the Stage

## » The landscape before we founded the company

- GPUs had 10x the flops of CPUs: nv40 vs. pentium 4
- Stanford had demonstrated the Brook project
- Lots of buzz about "GPGPU": What else can GPUs do?

#### » Brook

- What was Brook?
- Research developed in the Stanford Graphics Lab
  - » Pat Hanrahan, Ian Buck, Tim Foley, Daniel Horn, Jeremy Sugerman, Mike Houston, Kayvon Fatahalian
- Demonstrated HPC codes running on GPUs
  - Using compiler technology to make it work
- An open source project today

## **Many-Core Processors**

#### » There is a large category of Many-Core Processors

- GPUs: AMD & NVIDIA
- IBM Cell Processor
- Many-core CPUs: AMD & Intel
- Future: AMD Fusion Processor = CPU+GPU Integration

#### » Processor characteristics

- High memory bandwidth
- Extremely high flops
- High flop to memory access ratio
- On-chip communication network

## » Why use many-core processors?

- Performance
- Power
- Cost

## **Many-Core Processors**

- » Are many-core processors new?
  - No
- » Also called Stream Processors
  - Imagine Processor, Bill Daly et al, Stanford
  - Merrimack Architecture, Bill Daly et al, Stanford
  - SPI, Chief Scientist: Bill Daly
- » GPU architecture was heavily influenced by Stream Processors
  - As is the IBM Cell processor

#### Who Wants All These FLOPs?

## » Gaming

- Physics
- Image Processing
- AI? This has not yet been demonstrated, but it's intriguing

## » Image Processing

- Image & Video Editing
- Consumer & Professional

## » High Performance Computing

- Applications are solving big science problems numerically
- Server compute farms: from 1,000s to 100,000s of CPUs
- Workstations: CAD & Content. These have GPUs already
- Embedded: Medical & Defense

## What is High Performance Computing?

## » HPC uses computation to solve a science problem

- Oil & Gas: Seismic analysis, reservoir modeling...
- Finance: Monte carlo Simulations...
- Biology: Molecular modeling, sequence matching...
- Engineering: Fluid synamics...
- Government Labs: Stockpile simulation, climate...

## » Who are HPC Developers?

- Mostly scientists, but not computer scientists
- Mostly not parallel programming experts
- Mostly like programming in MatLab
- They are more interested in their science than in they are in optimizing a computer program

## What's Wrong with Multi-Core CPUs and GPUs?

## » Developer Productivity

- Most developers do not know how to write fast numerical codes
- Making x86 run fast is hard. GPUs are even harder.
- Developing threaded applications is hard (OpenMP & pthreads)
- Writing message-passing applications is very hard (MPI, Cell)

## » University curricula in numerical computing have shifted to high productivity languages

- MatLab: This is the tool of choice in hard sciences
- Scientists no longer learn Fortran
- Scientists are not computer scientists
- Scientists are not parallel programming experts
- Observation: MatLab is not a high performance system

## The PeakStream Programming Model

#### » We call it Stream Programming

- A data-parallel programming model
- With an explicit I/O model
- For many-core processors

#### » High performance

• The raison d'etre!

#### » Portable

- Across processor vendors, across processor generations
- (But does require significant effort by PeakStream)

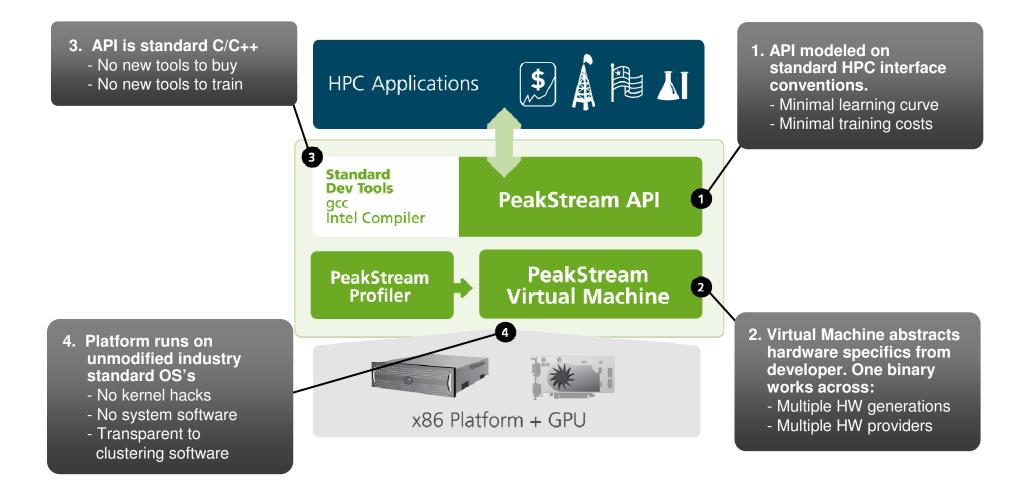
#### » Interoperable

Leverage existing libraries, tools, and systems (MPI, gcc, etc.)

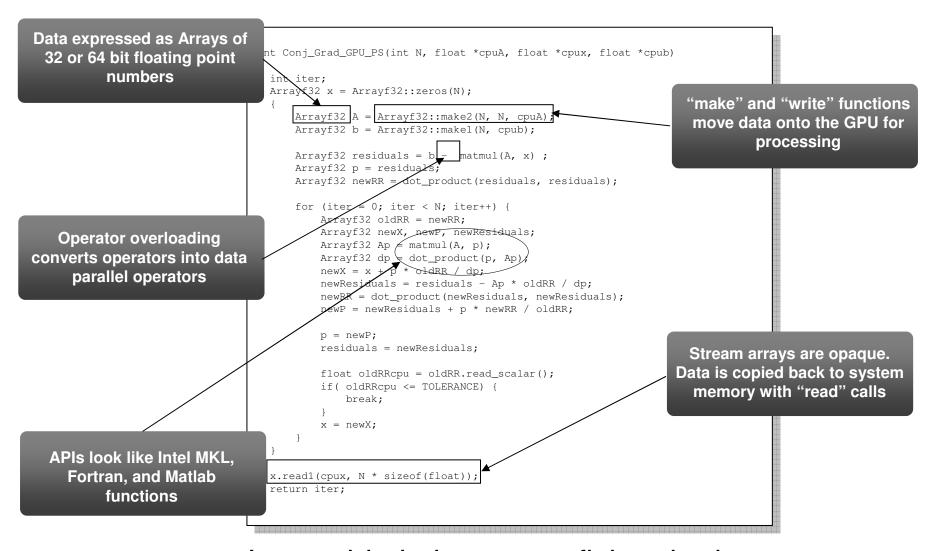
#### » High productivity

- Minimize time to solution
- For scientists & mathematicians
- Tools are important: debugger & profiler

## The PeakStream Platform™



## **PeakStream Programming Essentials**



Learnable in hours, proficient in days



## Why an API?

#### » New languages are rarely adopted

- They have steep learning curve
- They require new software ecosystems
  - » Compilers
  - » Tools
  - » Libraries
- Language extensions are new languages
  - » Definition of a new language: "won't compile with an existing compiler"

#### » APIs are much easier to adopt

- APIs are language-neutral
  - » They allow people to use their favorite languages
  - » They allow multiple language bindings: C, C++, Fortran, Java, ...
- APIs facilitate interoperability with existing software ecosystems
  - » MPI, OpenMP, MKL, ACML, ...
- APIs and languages are equally expressive

## **Virtual Machine with Dynamic Compilation**

## » Dynamic compilation facilitates binary portability

## » Across processor vendors

- Dynamically compile and optimize for the processor at hand
- NVIDIA and ATI GPUs have totally different ISAs
- GPUs and CPUs have very different ISAs and OS interfaces

## » Across processor generations

- Processors change faster than applications
- Want applications to automatically get faster as hardware gets faster
- But GPU ISAs change completely from one generation to the next
- Even x86 adds new instructions: SSE[1,2,3,4]

## **Dynamic Compilation Well Suited for HPC**

- » Dynamic compilation is now commonplace
  - GPU drivers
  - Java and .NET
  - Transmeta
  - VMware, XenSource
- » Dynamic compilation is fast
  - VMware running windows boots in 30 seconds
    - » Just 1 second of that is JIT code translation
- » Code caching is tremendously effective for HPC
  - Long running
  - Highly repetitive
- » JIT overhead easily amortized for HPC
  - High data-to-code ratio

## Computing $\pi$ with PeakStream

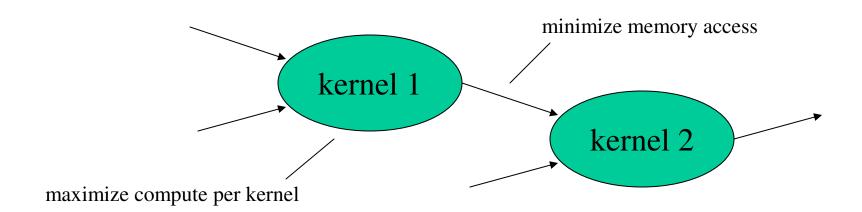
```
#include <peakstream.h>
#define NSET 1000000
                            // number of monte carlo trials
Arrayf32 Pi = compute_pi();
                            // get the answer as a 1x1 array
printf("Value of Pi = %f\n", pi);
Arrayf32
compute_pi(void)
  RNGf32 G(SP_RNG_DEFAULT, 271828); // create an RNG
  Arrayf32 X = rng\_uniform\_make(G, NSET, 1, 0.0, 1.0);
  Arrayf32 Y = rng_uniform_make(G, NSET, 1, 0.0, 1.0);
  Arrayf32 distance_from_zero = sqrt(X * X + Y * Y);
  Arrayf32 inside circle = (distance from zero <= 1.0f);</pre>
  return 4.0f * sum(inside circle) / NSET;
```

## What the Compiler Does: Generate Compute Kernels

- » Compiler outputs a series of Compute Kernels
  - And the VM executes them on the processor
- » A compute kernel is structured as:
  - Gather
  - Compute
  - Scatter
- » Maximize flops/kernel
  - Minimize memory bandwidth requirements
  - Avoid the processor memory wall
- » All of this is done automatically by the PeakStream JIT Compiler

## **The Stream Programming Paradigm**

- » Computation expressed as composition of compute *kernels*:
  - Gather phase
  - Compute phase
  - Scatter phase
- » Translates memory latency into memory bandwidth
  - Able to exploit processors with high compute/memory access ratios



## Computing $\pi$ with PeakStream

## » This is the code the VM generates and runs:

RNG & element-wise ops. reduction passes

final  $\pi$  calculation

## Detail of pass 1 GPU code:

S\_OUTPUT main(float2 THR\_ID : VPOS) {
 PS\_OUTPUT output;
 float4 tmp0, tmp1, tmp2, tmp3, tmp4,
 tmp5, tmp6, tmp7, tmp8, tmp9,
 tmp10;

tmp0 = CEICG12m6\_1d(in0, THR\_ID,
 inC0, inC1, inC2, inC3, inC4,
 inC5, out0\_pad);

tmp1 = smk32\_mul(tmp0, inC6.x);

tmp2 = smk32\_add(tmp1, inC7.x);

tmp3 = smk32\_mul(tmp2, tmp2);

tmp4 = CEICG12m6\_1d(in0, THR\_ID,
 inC8, inC9, inC10, inC11, inC12,
 inC13, out0\_pad);

tmp5 = smk32\_mul(tmp4, inC14.x);

tmp6 = smk32\_add(tmp5, inC15.x);

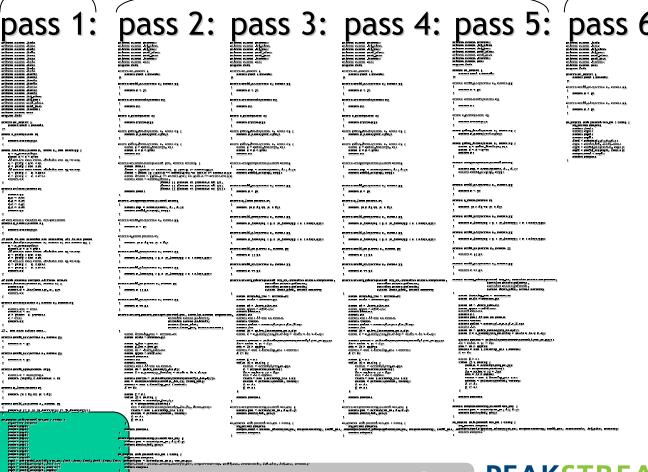
tmp7 = smk32\_mul(tmp6, tmp6);

tmp8 = smk32\_add(tmp3, tmp7);

tmp9 = smk32\_sqrt(tmp8);

tmp10 = smk32\_le(tmp9, inC16.x);

return output:



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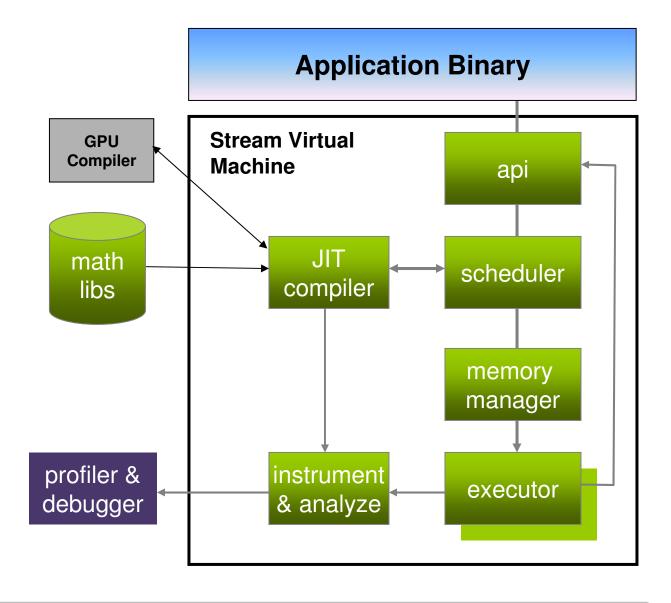
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## **Automatic Stream Kernel Synthesis**

## » Identifying the streaming kernel

- What's the granularity of the inner loop?
- How many GPU passes are optimal?
- » It's inappropriate for the application to pick
  - It is very processor-dependent
  - Depends on processor family, model, memory, ...
- » This is a good task for compilers
  - This is what the PeakStream JIT compiler does
  - Ensures portability of application code
  - Ensures scalable performance over many processors

## **PeakStream Software Architecture**



## **PeakStream Platform Functionality**

Standard Math  » Standard operators  » Range of logarithms  » Exp, powers, roots  » Rounding, abs etc.	Trigonometry  » Standard trigonometry  » Inverse trigonometry  » Hyperbolic functions	Array Manipulation  » Attribute queries  » Gather/Spread  » Indexing
Array Reduction & Statistics  » Sum/Min/Max  » Mean/Variance/Std Dev  » Random number generators	BLAS  » Full BLAS equivalance  » Levels 1/2/3	Signal processing  » Convolution  » Multiple border options  » 1D, 2D FFTs
Linear Algebra  » Dot product/Transpose  » Matmul  » LU & Cholesky Solvers	Array generation  » Identity and zero arrays  » Random number arrays  » Data stride	Utility functions  » Data transfer  » Performance hints  » VM management  » Debug APIs

Platforms & language	» Linux: Redhat Enterprise Linux, CentOS, gcc, gdb, icc
support	» Windows: WinXP SP2, Visual Studio, icc
	» C/C++

#### **Mandelbrot Fractal**

```
Arrayf32 Iter = -1;
Arrayf32 XP = (Arrayf32::index(0,pixels_x,pixels_y)-pixels_x/2);
Arrayf32 YP = (Arrayf32::index(1,pixels x,pixels y)-pixels y/2);
Arrayf32 Xprime = XP*cos(phi) - YP*sin(phi);
Arrayf32 Yprime = YP*cos(phi) + XP*sin(phi);
XP = Xprime/(pixels x*zoom)+cx;
YP = Yprime/(pixels_x*zoom)+cy;
for (int iteration=0; iteration<max iter; iteration++)</pre>
    // Iterate
    Arrayf32 Y=2*X*Y+YP;
    Arrayf32 X=X*X-Y2+XP;
    Arrayf32 Y2=Y*Y;
    // Test for escape condition
    Arrayf32 Eval = cond(X*X+Y2<4,0,1);
    Iter = cond(Iter<0&&Eval>0, iteration, Iter);
}
```

#### **PeakStream Linux Tools Extensions**

## **Debugger: gdb Extensions**

- » Debug PeakStream applications with a plug-in to the standard gdb debugger
- » Set breakpoints
- » Step through code executing on GPU & CPU
- » Examine arrays resident on the GPU
- » Generate reference results to compare GPU execution to CPU execution
- » Trap runtime errors

## **Profiler: Collection and Analysis**

- » Insight into optimization potentials
- » gprof style tool for analyzing application performance
- » Shows time spent per line and per function
- » Pinpoints excess data movement from system to local memory
- » View stream processor compute kernels
- » Analyze memory utilization

## PeakStream Debugger

- » GDB debugger extensions to monitor PeakStream arrays
- » Script provided for access

```
ps_gdb program
```

» DDE (Debugger Data Examination)

```
psprint array (print contents of SP array)
SP::DDE::get_array_element(A, idx0, idx1)
SP::DDE::read1(A, outptr, size, stride)
SP::DDE::read2(A, outptr, size, stride, pad)
SP::DDE::write array to file(A, filename)
```

#### » Error handlers

Either handle from your application or catch in the debugger

#### » Generate reference results

- To compare GPU to CPU results
- From your debugger session or your application

#### **PeakStream Profiler**

## » A gprof-style application profiler

- » Usage:
  - ps\_analyzer [options] [ > outfile ]

## » 3 basic views

- API Call view
- Compute kernel summary view
- Compute kernel detail view

## **Profiler: API Call View**

% Total Time	% Total	Calls	Compute	I/O	VM	API Name	Caller Name	File	Line
Cumulative	Time		Time	Time	Time				
29.48	29.48	100	0.075	0	0.032	<pre>SP::rng_uniform_make</pre>	GPU_PS_Compute_Pi	main.cpp	142
58.94	29.45	100	0.075	0	0.032	<pre>SP::rng_uniform_make</pre>	GPU_PS_Compute_Pi	main.cpp	143
85.27	26.33	10	0.037	0	0.059	SP::sum	GPU_PS_Compute_Pi	main.cpp	146
90.75	5.48	10	0.014	0	0.0059	SP::operator<=	GPU_PS_Compute_Pi	main.cpp	145
95.54	4.79	10	0.012	0	0.0052	SP::sqrt	GPU_PS_Compute_Pi	main.cpp	144
96.35	0.81	10	0	0.0029	0	SP::RNGf32::RNGf32	GPU_PS_Compute_Pi	main.cpp	140
97.06	0.71	10	0.00013	0	0.0025	SP::operator/	GPU_PS_Compute_Pi	main.cpp	146
97.75	0.68	10	0.0017	0	0.00074	SP::cond	GPU_PS_Compute_Pi	main.cpp	145
98.43	0.68	10	0.0017	0	0.00074	SP::operator+	GPU_PS_Compute_Pi	main.cpp	144
99.12	0.68	10	0.0017	0	0.00074	SP::operator*	GPU_PS_Compute_Pi	main.cpp	144
99.80	0.68	10	0.0017	0	0.00074	SP::operator*	GPU_PS_Compute_Pi	main.cpp	144
99.94	0.14	10	2.7e-05	0	0.00049	SP::operator*	GPU_PS_Compute_Pi	main.cpp	146
100.00	0.06	10	0	0.0002	0	<pre>SP::Arrayf32::read_scalar</pre>	GPU_PS_Compute_Pi	main.cpp	148

#### » Role: report on how the application spent its time

in terms of PeakStream API calls

## » Conclusions for this simple example:

- This application is not I/O limited
- Most of the run time is spent in the RNG & reduction

## **Profiler: Compute Kernel Summary View**

% Total	Executions	Compute	Paging	JIT	Kernel	File	Line
Time		Time	Time	Time	Name		
71.92	10	0.18	0	0.078	GPU_PS_Compute_Pi:1	main.cpp	142
26.33	60	0.037	0	0.059	GPU_PS_Compute_Pi:2	main.cpp	146
0.85	10	0.00016	0	0.0029	GPU_PS_Compute_Pi:3	main.cpp	146

» Role: report on which compute kernels matter most

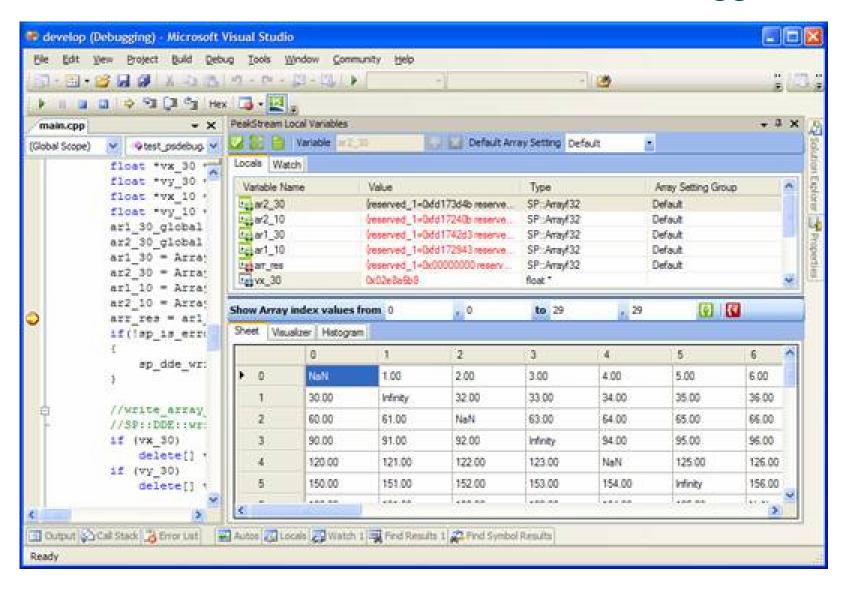
## **Profiler: Compute Kernel Detail View**

% Total Time	Execu		Compute Time	Paging Time	JIT Time	Kernel Name		File		Line
71.92		10	0.18	0	0.078	GPU_PS_C	Compute_Pi:1	main.cp	p	142
	Detail	s:								
		Comput	e JIT	API	Name		Caller		File	Line
		Time	Time	<b>e</b>						
	0:	0.07	2 0.03	3 SP:	:rng_unifo	cm_make	GPU_PS_Co	ompute_Pi	main.cpp	142
	1:	0.001	7 0.0007	4 SP:	:rng_unifor	rm_make	GPU_PS_Co	mpute_Pi	main.cpp	142
	2:	0.001	7 0.0007	4 SP:	:rng_unifor	rm_make	GPU_PS_Co	mpute_Pi	main.cpp	142
	3:	0.001	7 0.0007	4 SP:	:operator*		GPU_PS_Co	mpute_Pi	main.cpp	144
	4:	0.07	2 0.03	3 SP:	:rng_unifor	rm_make	GPU_PS_Co	mpute_Pi	main.cpp	143
	5:	0.001	7 0.0007	4 SP:	:rng_unifor	rm_make	GPU_PS_Co	mpute_Pi	main.cpp	143
	6:	0.001	7 0.0007	4 SP:	:rng_unifor	rm_make	GPU_PS_Co	mpute_Pi	main.cpp	143
	7:	0.001	7 0.0007	4 SP:	:operator*		GPU_PS_Co	mpute_Pi	main.cpp	144
	8:	0.001	7 0.0007	4 SP:	:operator+		GPU_PS_Co	mpute_Pi	main.cpp	144
	9:	0.01	2 0.0052	2 SP:	:sqrt		GPU_PS_Co	mpute_Pi	main.cpp	144
	10:	0.01	4 0.005	9 SP:	:operator<	=	GPU_PS_Co	mpute_Pi	main.cpp	145
	11:	0.001	7 0.0007	4 SP:	:cond		GPU_PS_Co	mpute_Pi	main.cpp	145

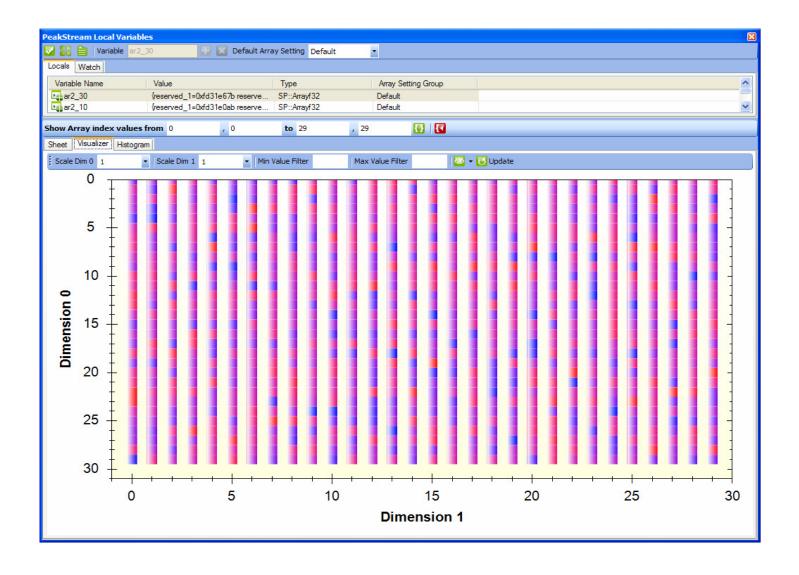
» Role: report about what's inside a compute kernel



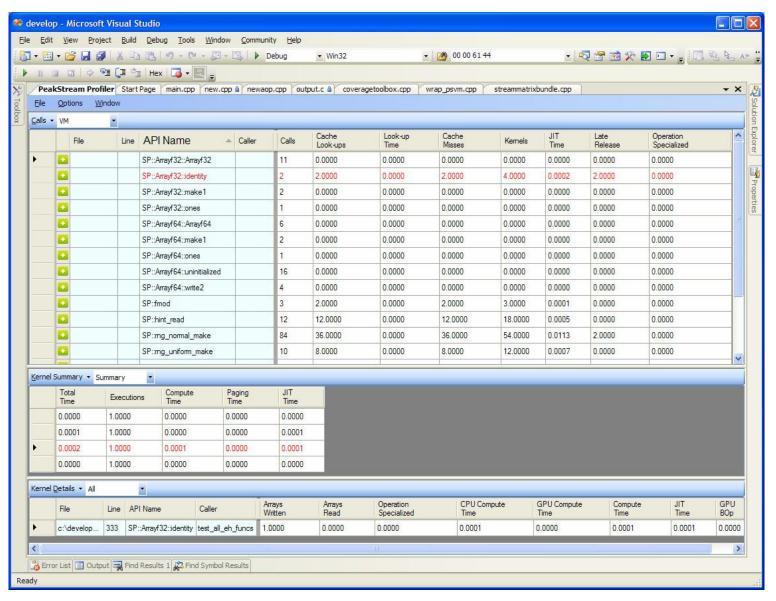
## PeakStream Tools Extensions: Windows Debugger



## **PeakStream Tools Extensions: Windows Visualizer**

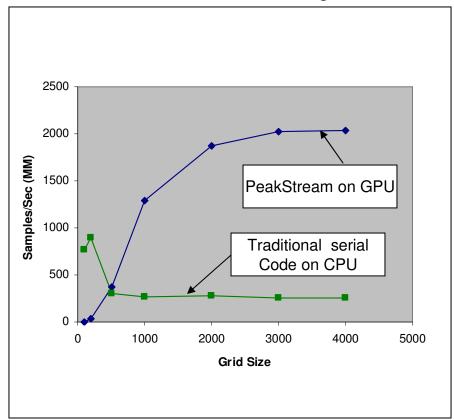


## PeakStream Tools Extensions: Windows Profiler

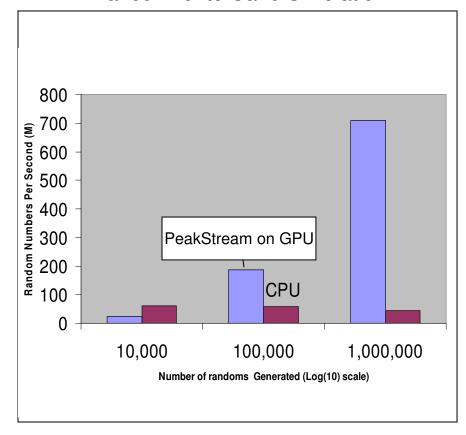


## **Lab Application Benchmarks**

Oil & Gas: Kirchhoff Migration



Finance: Monte Carlo Simulation



8x Peak Performance Advantage

16x Peak Performance Advantage

## **Application: Kirchhoff Migration**

```
void
KirchhoffMigration(int NT, int N, float *datagpu, float *modlgpu)
 int NTN = NT * N;
  float dx = LX / float(N);
  float dt = LT / float(NT);
  float factor = 1./ (velhalf * velhalf);
  float idt = 1./ dt;
 Arrayf32 mod1 = zeros_f32(NT,N);
    Arrayf32 x = dx * index_f32(1, NT, N);
    Arrayf32 z = dt * index_f32(0, NT, N);
    Arrayf32 data = Arrayf32::make2(NT, N, datagpu);
    for(int iy=0; iy < N; iy++) {</pre>
      float y = float(iy)*dx;
      Arrayf32 index1 = float(iy) * ones_f32(NT, N);
      Arrayf32 it = 0.5 + \text{sqrt}(z * z + (x-y)* (x-y) * \text{factor}) * \text{idt};
      modl += gather2_floor(data, it, index1);
    }
  modl.read1(modlgpu, NTN * sizeof(float) );
  return;
}
```

## **Application: Monte Carlo Finance**

```
float MonteCarloAntithetic(float price, float strike, float vol,
                                        float rate, float div, float T )
 float deltat
                    = T/N:
 float muDeltat
                    = (rate-div-0.5*vol*vol)*deltat;
 float volSqrtDeltat = vol*sqrt(deltat);
 float meanCPU
                     = 0.0f;
 Arrayf32 meanSP; // result
                          // a new scope to hold temporary arrays
   RNGf32 rng hndl(SP_RNG_CEICG12M6, 0);
   Arrayf32 U = zeros_f32( M );
   for(int i=0; i<N; i ++) {
     U += rnq normal make(rnq hndl, M);
   Arrayf32 values;
     Arrayf32 lnS1 = log(price) + N * muDeltat + volSqrtDeltat*U;
     Arrayf32 lnS2 = log(price) + N * muDeltat + volSqrtDeltat*(-U);
     Arrayf32 S1 = exp(lnS1);
     Arrayf32 S2 = exp(lnS2);
     values = (0.5 * (max(0,S1-strike) + max(0,S2-strike)) * exp(-rate*T));
   meanSP = mean( values );
                          // all temporaries released as we exit scope
 meanCPU = meanSP.read_scalar();
  return meanCPU;
```

#### The Future of Processors

- » Where are Processors going?
- » Integrated CPU + GPU
  - AMD's Fusion
  - Intel's Larrabee
- » GPUs and Cell Processor are not so different
  - They will converge
  - And become integrated with a few CPU cores
  - What differentiates them from CPUs? Explicit communication models
- » The Future Processor
  - A control processor (a.k.a. CPU)
  - A compute array (GPU, Cell, etc.)
- » This is an excellent processor for both gaming and HPC

## Where is Software for Many-Core Processors Going?

## » New programming models

- Data-parallel is one approach, and there are others
- But the manually threaded approach leaves a lot to be desired
- How do we expose new models? APIs & languages

#### » Increasing importance of runtime systems

- Managing the processors: Scheduling
- Managing the data: Memory Management
- Managing the code: JITing

## » Reliance on compilers

- To create optimal compute kernels for rapidly evolving processors
- In a way that protects the investment in application codes

#### Conclusion

- » The world needs good programming environments to make parallel programming easier
  - This is an exciting area of continued research
  - The need will persist for a long time
- » PeakStream was one such solution
  - A data-parallel model for programming many-core
  - What other solutions can you think of?
- » Thank you very much