The Power of Abstraction

Barbara Liskov October 2010

Outline

- Inventing abstract data types
- CLU
- Type hierarchy
- What next

Data Abstraction Prehistory

The Venus machine

The Interdata 3

- par	
Model 3	INTERDATA
Q	
-Q	
0	

Data Abstraction Prehistory

The Venus machineThe Venus operating system

Data Abstraction Prehistory

- The Venus machine
- The Venus operating system
- Programming methodology

Programming Methodology

- How should programs be designed?
- How should programs be structured?

 E. W. Dijkstra. Go To Statement Considered Harmful. Cacm, Mar. 1968

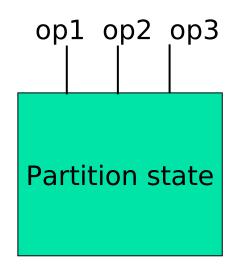
 N. Wirth. Program Development by Stepwise Refinement. Cacm, April 1971

- D. L. Parnas. Information Distribution Aspects of Design Methodology. IFIP Congress, 1971
- "The connections between modules are the assumptions which the modules make about each other."

Partitions

B. Liskov. A Design Methodology for Reliable Software Systems. FJCC, Dec. 1972





From Partitions to ADTs

How can these ideas be applied to building programs?



Connect partitions to data types

Meeting in Savanah

- ACM Sigplan-Sigops interface meeting. April 1973. (Sigplan Notices, Sept. 1973)
- Started to work with Steve Zilles

Extensible Languages

- S. Schuman and P. Jourrand.
 Definition Mechanisms in Extensible Programming Languages. AFIPS. 1967
- R. Balzer. Dataless Programming.
 FJCC 1967

 O-J. Dahl and C.A.R. Hoare. Hierarchical Program Structures. Structured Programming, Academic Press, 1972

J. H. Morris. Protection in Programming Languages. Cacm. Jan. 1973

 W. Wulf and M. Shaw. Global Variable Considered Harmful. Sigplan Notices. Feb. 1973.

Abstract Data Types

 B. Liskov and S. Zilles. Programming with Abstract Data Types. ACM Sigplan Conference on Very High Level Languages. April 1974

What that paper proposed

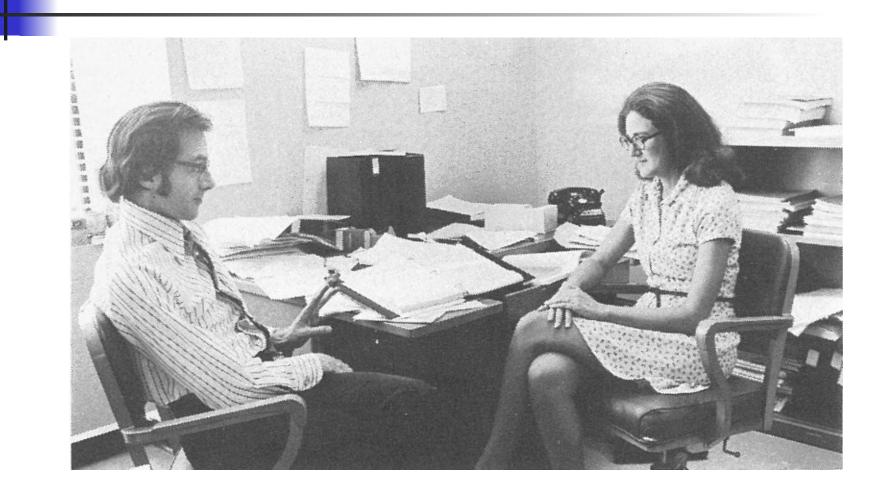
- Abstract data types
 - A set of operations
 - And a set of objects
 - The operations provide the only way to use the objects

What that paper proposed

- Abstract data types
 - Clusters with encapsulation
- Polymorphism
- Static type checking (we hoped)
- Exception handling

From ADTs to CLU

- Participants
 - Russ Atkinson
 - Craig Schaffert
 - Alan Snyder



Why a Programming Language?

- Communicating to programmers
- Do ADTs work in practice?
- Getting a precise definition
- Achieving reasonable performance

Language Design

- Goals
 - Expressive power, simplicity, performance, ease of use
 - Minimality
 - Uniformity
 - Safety

Language Design

- Restrictions
 - No concurrency
 - No go tos
 - No inheritance

Some Assumptions/Decisions

- Heap-based with garbage collection!
- No block structure!
- Separate compilation
- Static type checking

CLU Mechanisms

- Clusters
- Polymorphism
- Exception handling
- Iterators

Clusters

IntSet = cluster is create, insert, delete, isln, ...

end IntSet

Clusters

IntSet = cluster is create, insert, delete, ... end IntSet

IntSet s := IntSet\$create()
IntSet\$insert(s, 3)



IntSet = cluster is create, insert, delete,

rep = array[int]

Clusters

IntSet = cluster is create, insert, delete,

rep = array[int]

create = proc () returns (cvt)
 return (rep\$create())
 end create

Polymorphism

Set = cluster[T: type] is create, insert,
...
end Set

Set[int] s := Set[int]\$create()
Set[int]\$insert(s, 3)

Polymorphism

Set = cluster[T: type] is create, insert, ...
where T has equal: proctype(T, T)
returns (bool)

Polymorphism

Set = cluster[T: type] is create, insert, ...
where T has equal: proctype(T, T)
returns (bool)

rep = array[T]

insert = proc (x: cvt, e: T) ... if e = x[i] then ...

- J. Goodenough. Exception Handling: Issues and a Proposed Notation. Cacm, Dec. 1975
 - Termination vs. resumption
 - How to specify handlers

. . .

choose = proc (x: cvt) returns (T)
signals (empty)
if rep\$size() = 0 then signal empty

choose = proc (x: cvt) returns (T)
signals (empty)
if rep\$size() = 0 then signal empty

set[T]\$ choose(s)
 except when empty: ...

. . .

- Handling
- Propagating
- Shouldn't happen
 - The failure exception
- Principles
 - Accurate interfaces
 - Avoid useless code



For all x in C do S

Iterators

For all x in C do S

- Destroy the collection?
- Complicate the abstraction?

Visit to CMU

- Bill Wulf and Mary Shaw, Alphard
- Generators



sum: int := 0
for e: int in Set[int]\$members(s) do
 sum := sum + e
 end

Iterators

Set = cluster[T] is create, ..., members, ...

rep = array[T]

members = iter (x: cvt) yields (T)
for z: T in rep\$elements(x) do
 yield (z) end



Argus and distributed computingType Hierarchy

The Landscape

- Inheritance was used for:
 - Implementation
 - Type hierarchy

Type hierarchy

- Wasn't well understood
- E.g., stacks vs. queues

The Liskov Substitution Principle (LSP)

 Objects of subtypes should behave like those of supertypes if used via supertype methods

 B. Liskov. Data abstraction and hierarchy. Sigplan notices, May 1988



 Modularity based on abstraction is the way things are done

Challenges

- New abstraction mechanisms?
- Massively Parallel Computers
- Internet Computer
 - Storage and computation
 - Semantics, reliability, availability, security

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