CME307/MS&E311 Theory Review

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Optimization Problems

- A set of decision variables, *x*, in vector or matrix form with dimension n
- A continuous and sometime differentiable objective function *f(x)*
- A feasible region where *x* can be in

min	$f(\boldsymbol{x})$
s.t.	$\boldsymbol{x} \in \boldsymbol{X}$

- One can smooth them by reformulation as constrained optimization:

$$\max \quad \min_i\{ f_i(x), i=1,\ldots,n\} \;\text{->}$$

$$\max \quad \alpha \qquad \text{s.t. } \alpha - f_i(x) \le 0, \text{ for } i=1,\ldots,n$$

Function, Gradient Vector and Hessian Matrix

- A function f of x in R^n
- The Gradient Vector of f at x

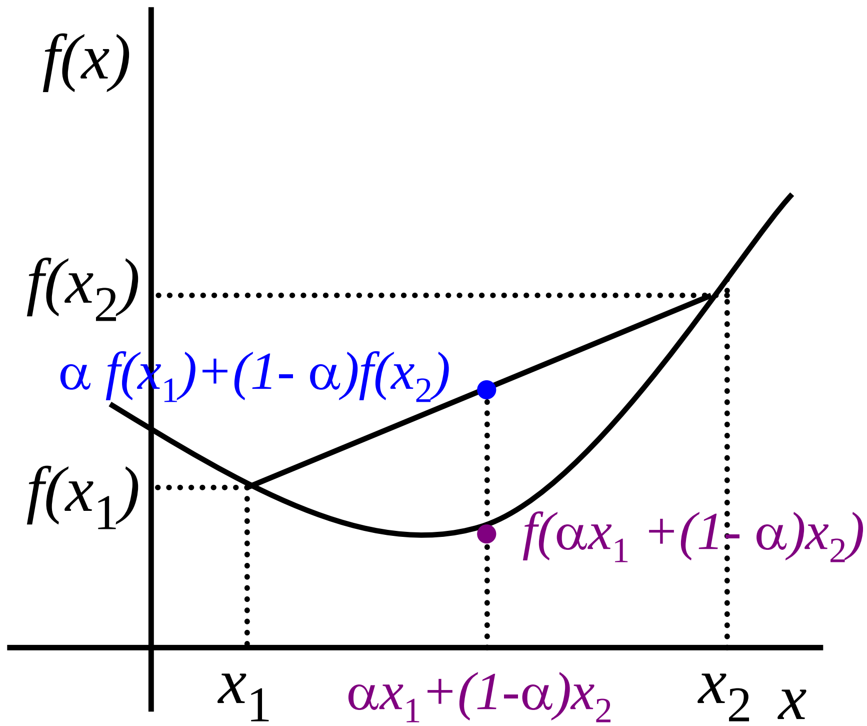
$$\nabla f(x) = \left(\frac{\partial f}{\partial x_1} \quad \frac{\partial f}{\partial x_2} \cdots\cdots\cdots \frac{\partial f}{\partial x_n} \right)$$

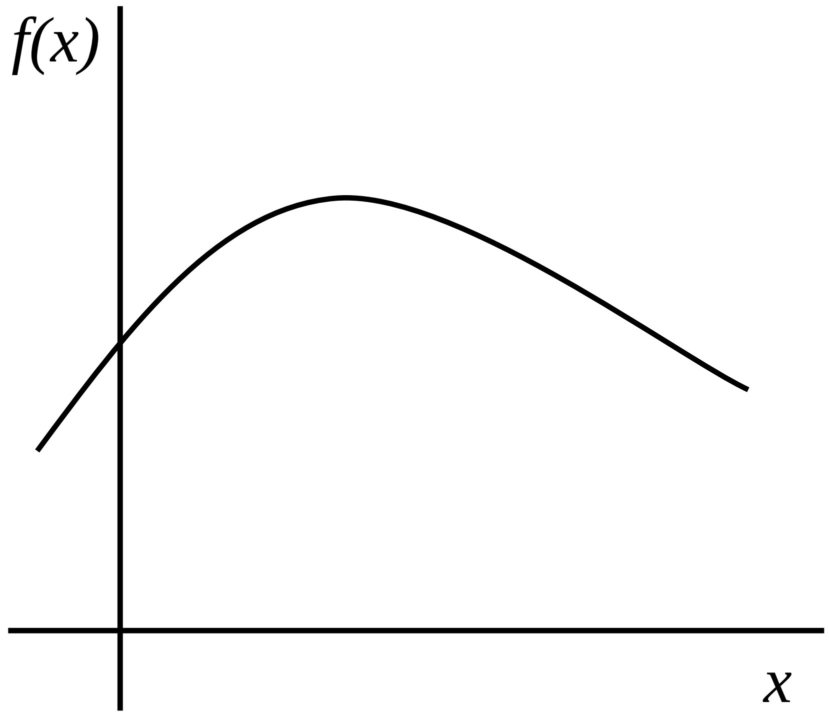
- The Hessian Matrix of f at x

$$\nabla^2 f(x) = \begin{pmatrix} \frac{\partial f^2}{\partial x_1 \partial x_1} & \cdots\cdots\cdots & \frac{\partial f^2}{\partial x_1 \partial x_n} \\ & \cdots\cdots\cdots\cdots & \\ \frac{\partial f^2}{\partial x_n \partial x_1} & \cdots\cdots\cdots\cdots & \frac{\partial f^2}{\partial x_n \partial x_n} \end{pmatrix}$$

- **Taylor's Expansion Theorem**

Convex and Concave Functions





f(x) is a <u>convex function</u> if and only if for any given two points x_1 and x_2 in the function domain and for any constant $0 \leq \alpha \leq 1$

$$f(\alpha x_1 + (1-\alpha)x_2) \leq \alpha f(x_1) + (1-\alpha) f(x_2)$$

Strictly convex if $x_1 \neq x_2$, $f(0.5x_1 + 0.5x_2) < 0.5f(x_1) + 0.5f(x_2)$

Convex Quadratic Functions

$f(x)=x^TQx+c^Tx$ is a convex function if and only if Hessian matrix Q is positive semi-definite (PSD).

$f(x)=x^TQx+c^Tx$ is a strictly convex function if and only if Q is positive definite (PD).

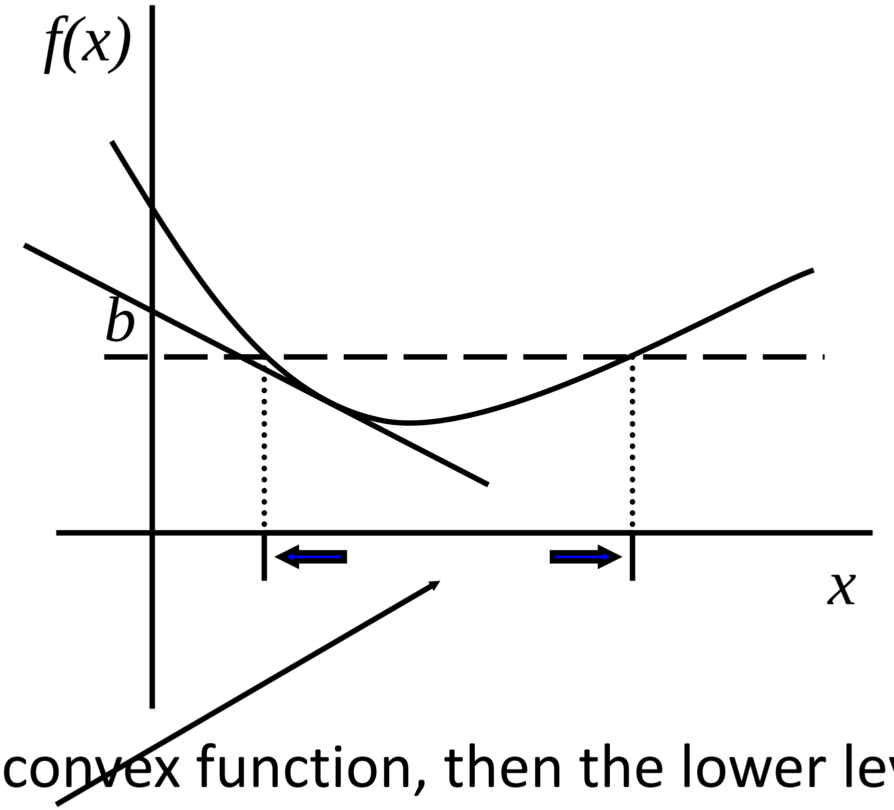
Q is PSD if and only if $x^TQx \geq 0$ for all x.

A 2x2 matrix is PSD (or PD) if and only if two diagonal entries and the determinant are nonnegative (or positive)

Convex Sets

- A set is convex if every line segment connecting any two points in the set is contained entirely within the set
  - Ex - polyhedron
  - Ex - ball
- An extreme point of a convex set is any point that is not on any line segment connecting any other two distinct points of the set
- The intersection of convex sets is a convex set
- A set is closed if the limit of any convergent sequence of the set belongs to the set

Properties of Convex Function



If *f(x)* is a convex function, then the lower level set $\{x: f(x) \le b\}$ is a convex set for any constant *b*.

The graph of a convex function lies above its <u>tangent line (planes)</u>.
The Hessian matrix of a convex function is <u>positive semi-definite</u>.

Optimization Problem Classes

- <u>Unconstrained Optimization</u>
  - Convex or Nonconvex
- <u>Constrained Optimization</u>
  - Conic Linear Optimization (CLO)
  - Convex Constrained Optimization (CCO)
    - Feasible region/set convex; objective general
  - Generally Constrained Optimization (GCO)
  - Convex Optimization (CO)
    - Minimize a convex function over a convex feasible set
    - Maximize a concave function over a convex feasible set
    - CLO belongs to CO

$$\begin{array}{ll} \min & f(\boldsymbol{x}) \\ \text{s.t.} & \boldsymbol{x} \in \boldsymbol{X} \end{array}$$

Optimization Problem Forms

$$\begin{array}{ll} \min & \boldsymbol{c}^T\boldsymbol{x} \\ \text{s.t.} & A\boldsymbol{x} - \boldsymbol{b} = 0, \\ & \boldsymbol{x} \in \boldsymbol{K} \end{array}$$

Conic Linear Optimization (CLO)

A: an m x n matrix
***c*: objective coefficient**
***K*: a closed convex cone**

This is convex optimization

$$\begin{array}{ll} \min & f(\boldsymbol{x}) \\ \text{s.t.} & c_i(\boldsymbol{x}) \le, =, \ge 0, \ i=1,\ldots,m \end{array}$$

Generally Constrained Optimization (GCO)

Each function can be continuous, continuously differentiable (C^1), or twice continuously differentiable (C^2)

It is Convex Optimization if c_i with "≤" are all convex functions, c_i with "≥" are all concave functions, c_i with "=" are all linear/affine functions, and *f* is a convex function

Why do we care about convex optimization?

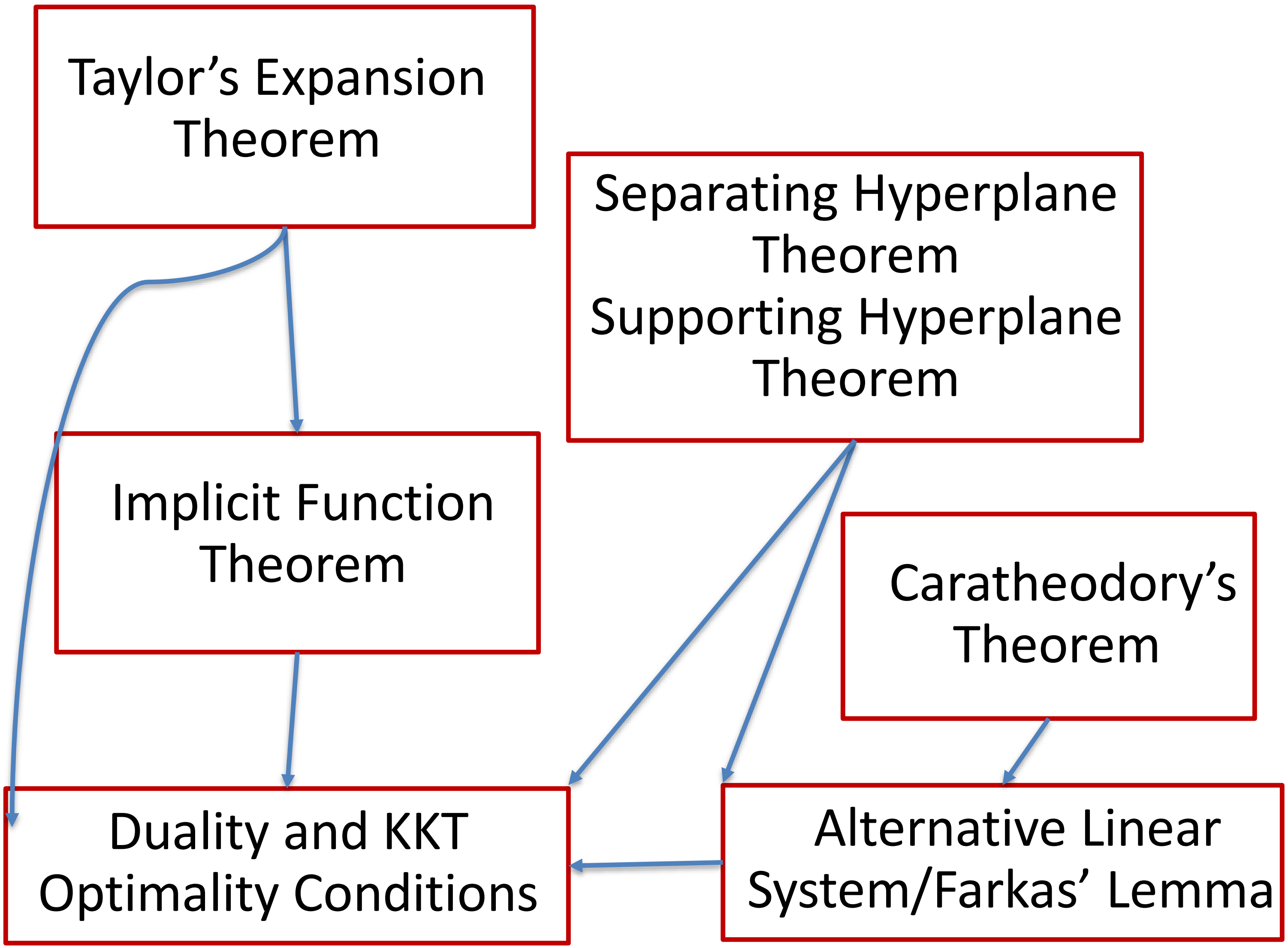
- It guarantees that every local optimizer is a global optimizer
- It guarantees that every (first-order) KKT (or stationary) point/solution is a global optimizer
- This is significant because all of our numerical optimization algorithms search/generate a KKT point/solution
- Sometime the problem can be “convexfied”:

$$\min \quad c^T x, \quad \text{s.t.} \quad ||x||^2=1$$

$$\Updownarrow$$

$$\min \quad c^T x, \quad \text{s.t.} \quad ||x||^2 \leq 1$$

Optimization **Theory**: Mathematical Foundations



Theory: Feasibility Conditions

- Feasibility Conditions or Farkas' Lemmas are developed to characterize and certify feasibility or infeasibility of a feasible region
- Alternative Systems X and Y: X has a feasible solution if and only if Y has no feasible solution
  - X and Y cannot both have feasible solution
  - Exactly one of them has a feasible solution
- They can be viewed as special cases of Linear Programming primal and dual pairs

Alternative Systems and CLO Pairs I

$$\begin{aligned} A\mathbf{x} - \mathbf{b} &= \mathbf{0}, \\ \mathbf{x} &\in K \end{aligned}$$

System X

A: an m x n matrix

***b*: m-dimension vector**

***K*: a closed convex cone**

$$\begin{aligned} \mathbf{b}^T\mathbf{y} &= 1 (>0) \\ A^T\mathbf{y} + \mathbf{s} &= \mathbf{0}, \\ \mathbf{s} &\in K^* \end{aligned}$$

System Y

K* **is the dual cone**

$$\begin{aligned} p^* = \min \quad & \mathbf{0}^T\mathbf{x} \\ \text{s.t.} \quad & A\mathbf{x} - \mathbf{b} = \mathbf{0}, \\ & \mathbf{x} \in K \end{aligned}$$

$$\begin{aligned} d^* = \max \quad & \mathbf{b}^T\mathbf{y} \\ \text{s.t.} \quad & A^T\mathbf{y} + \mathbf{s} = \mathbf{0}, \\ & \mathbf{s} \in K^* \end{aligned}$$

Alternative Systems and CLO Pairs II

$$\boldsymbol{c}^T\boldsymbol{x}=-1(<0)$$

$$A\boldsymbol{x} = 0,$$

$$\boldsymbol{x} \in K$$

$$A^T\boldsymbol{y} + \boldsymbol{s} - \boldsymbol{c} = 0,$$

$$\boldsymbol{s} \in K^*$$

System X
A: an m x n matrix
***c*: n-dimension vector**
***K*: a closed convex cone**

System Y

***K** is the dual cone**

p*=min $\boldsymbol{c}^T\boldsymbol{x}$

s.t. $A\boldsymbol{x} = 0,$

$\boldsymbol{x} \in K^*$

d*=max $0^T\boldsymbol{y}$

s.t. $A^T\boldsymbol{y} + \boldsymbol{s} - \boldsymbol{c} = 0,$

$\boldsymbol{s} \in K$

Feasibility Test Machine

Is system **X** feasible?

Yes → Is system **Y** feasible?

"Not" under any circumstances

No → Is system **Y** feasible?

"Yes" under certain conditions of cone *K* and data matrix *A*:

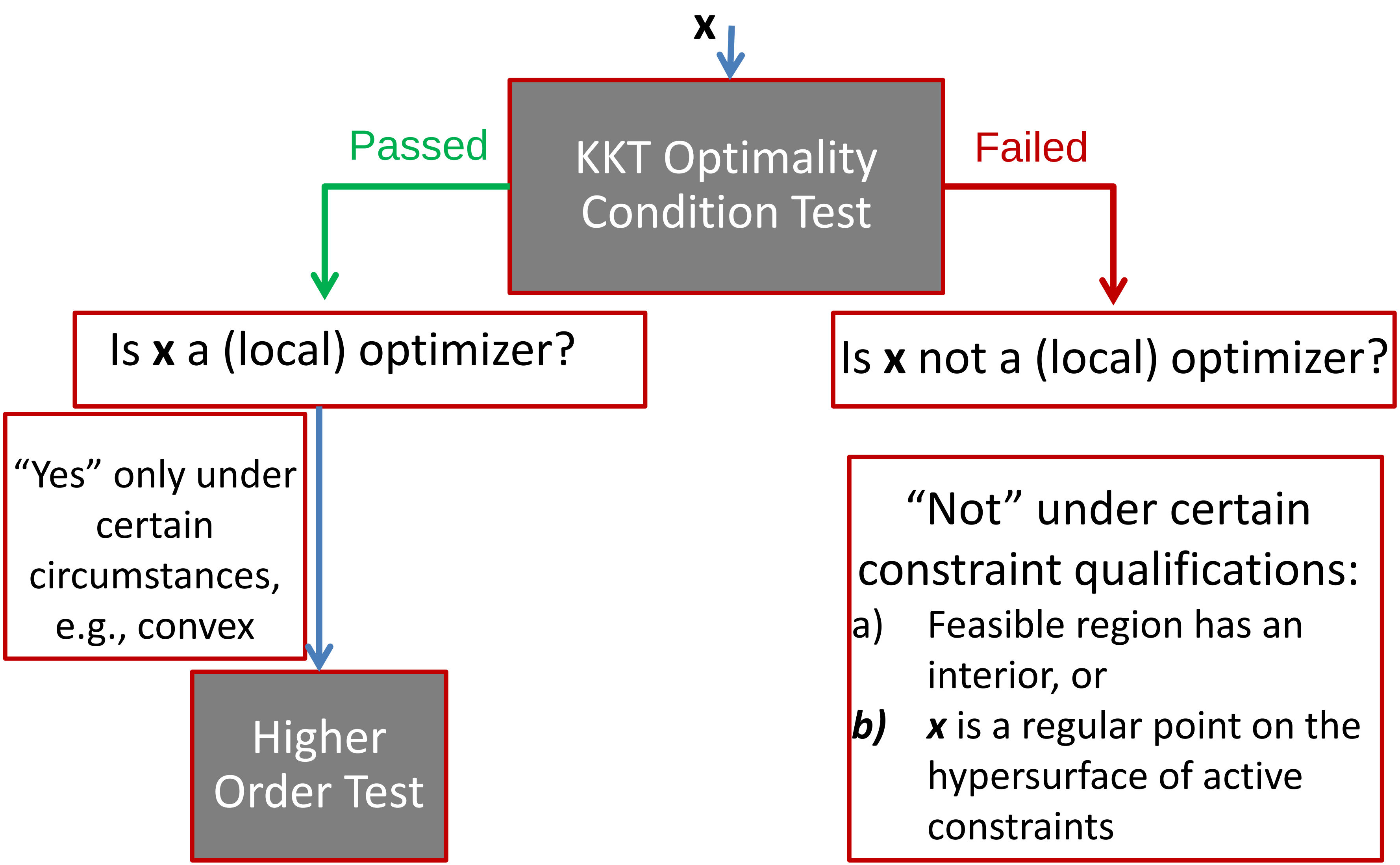
a) *K* is a polyhedron cone, or

b) $A\mathbf{x}$ *or* $A^T\mathbf{y}$ has an interior solution

Theory: Optimality Conditions

- Optimality (KKT) and Duality Conditions are developed to characterize and certify possible minimizers
  - Feasibility of original variables
  - Optimality conditions consist of original variables and Lagrange multipliers
  - Zero-order, First-order, Second-order, necessary, sufficient
- They may not lead directly to a very efficient algorithm for solving problems, but they do have a number of benefits:
  - They give insight into what optimal solutions look like
  - They provide a way to set up and solve small problems
  - They provide a method to check solutions to large problems
  - The Lagrange multipliers can be seen as sensitivities of the constraints
- A minimizers may not satisfy optimality conditions unless certain *constraint qualifications* hold.

KKT Optimality Condition Test Machine



Conic Duality Theorems for CLO

p^*=min	$\boldsymbol{c}^T\boldsymbol{x}$	d^*=max	$\boldsymbol{b}^T\boldsymbol{y}$
s.t.	$A\boldsymbol{x} - \boldsymbol{b} = \boldsymbol{0},$	s.t.	$A^T\boldsymbol{y} + \boldsymbol{s} - \boldsymbol{c} = \boldsymbol{0},$
	$\boldsymbol{x} \in \boldsymbol{K}$		$\boldsymbol{s} \in \boldsymbol{K}^*$

Primal Problem
A: an m x n matrix
***c*: objective coefficient**
***K*: a closed convex cone**

Weak Duality Theorem

Dual Problem

***K** is the dual cone**

0-Order Condition: $p^* = d^*$

Sufficient!

Strong Duality Theorem: They must equal?

"Yes" under certain conditions of cone K and data matrix A,b,c:

a) K is a polyhedron cone, or
b) *either one* has an interior feasible solution

The Lagrange Function of GCO

min	$f(\boldsymbol{x})$
s.t.	$c_i(\boldsymbol{x})\ (\le, =, \ge)\ 0$, i=1,…,m

Restriction on multipliers y_i,

y_i (≤,"free",≥) 0, $i=1,\dots,m$

The Largrange Function

$$L(\mathbf{x},\mathbf{y}) = f(\boldsymbol{x}) - \sum_i y_i c_i(\boldsymbol{x})$$

The Lagrange function can be interpreted as a "penalized" aggregated objective function:

y_i free: can be penalized either way
$y_i \ge 0$: can be penalized when $c_i(\boldsymbol{x}) \le 0$
$y_i \le 0$: can be penalized when $c_i(\boldsymbol{x}) \ge 0$

y: Sensitivity and Shadow Price interpretation for optimal-value

The Lagrangian Duality for GCO

p^*=min $f(\boldsymbol{x})$

s.t. $c_i(\boldsymbol{x})$ (≥,=,≤) 0, i=1,...,m

Weak Duality Theorem
$P^* \geq d^*$

Let $\phi(\boldsymbol{y}) = \inf_{\boldsymbol{x}} L(\boldsymbol{x},\boldsymbol{y})$

Strong Duality Theorem
They must equal?

d^*=max $\phi(\boldsymbol{y})$

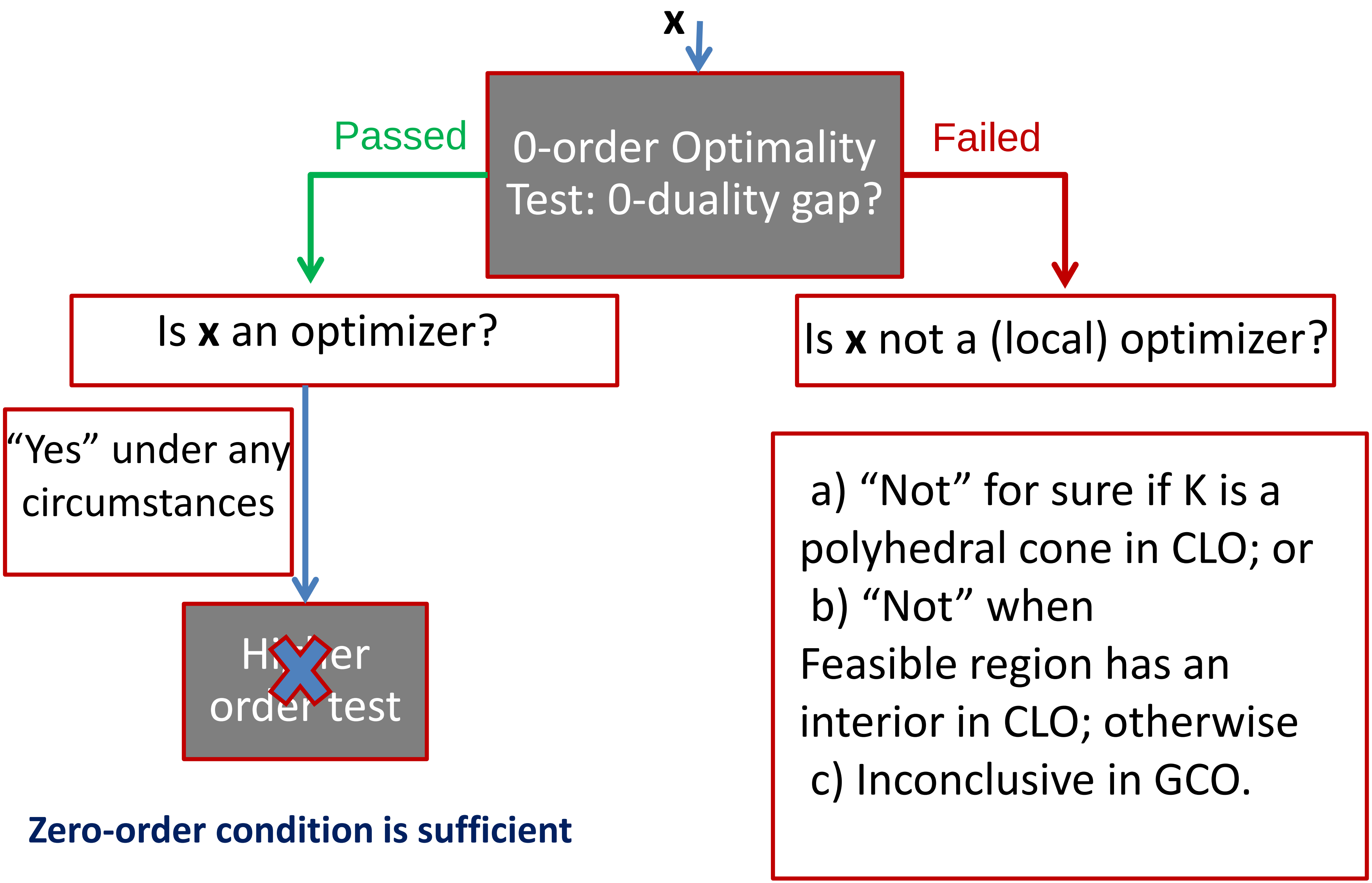
s.t. y_i (≤,"free", ≥) 0, $i=1,\ldots,m$

Not necessarily!

0-Order Condition: $p^* = d^*$
(to replace CSC)

Sufficient!

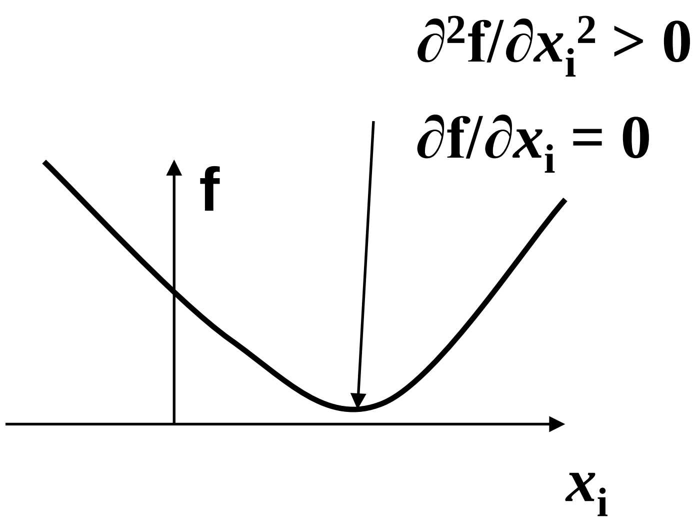
Zero-Order Optimality Test for CLO and GCO



1 and 2-order Conditions: Unconstrained

- Problem:
  - **Minimize *f(x)*, where *x* is a vector that could have any values, positive or negative**
- <u>First Order Necessary Condition (min or max)</u>:
  - **$\nabla f(x) = 0$ ($\partial f/\partial x_i = 0$ for all i) is the first order necessary condition for optimization**
- <u>Second Order Necessary Condition</u>:
  - **$\nabla^2 f(x)$ is positive semidefinite (PSD)**
    - **[$d^T \nabla^2 f(x) d \geq 0$ for all d]**
- <u>Second Order Sufficient Condition</u>

  (Given FONC satisfied)
  - **$\nabla^2 f(x)$ is positive definite (PD)**
    - **[$d^T \nabla^2 f(x) d > 0$ for all d ≠ 0]**



1-Order KKT Condition for GCO

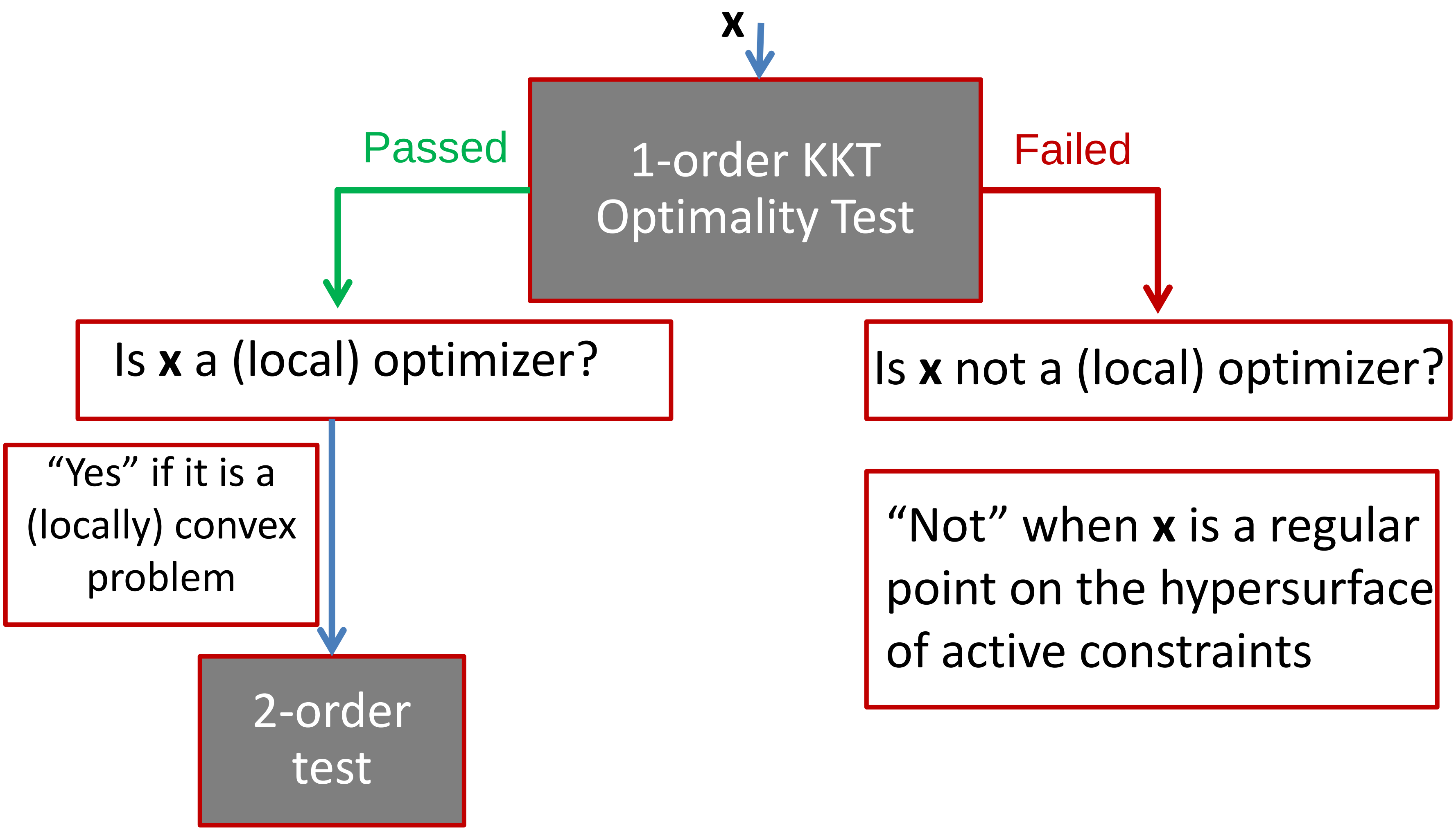
Recall the Largrange Function

$$L(\mathbf{x},\mathbf{y}) = f(\boldsymbol{x}) - \sum_i c_i(\mathbf{x}) y_i$$

$\nabla_x L(\mathbf{x},\mathbf{y}) = \mathbf{0}$, that is,
$\partial L(\mathbf{x},\mathbf{y})/\partial x_j = 0$, for all $j=1,\ldots,n$, and
$c_i(\mathbf{x}) y_i = 0$, for all $i=1,\ldots,m$
y_i (≤,"free", ≥) 0, $c_i(\boldsymbol{x})$ (≤, =, ≥) 0, $i=1,\ldots,m$

LDC
CSC
MSC and OPC

Optimality Test for GCO



Rules to Construct the Dual from 1-st Order KKT for CO

Primal or Original Problem Constraints:

$c_i(\boldsymbol{x})$ $(\geq,=,\leq)$ 0, i=1,…,m

Dual Constraints I:

y_i (≥,"free",≤) 0, $i=1,\ldots,m$

Dual Constraints II:

$\partial L(\mathbf{x},\mathbf{y})/\partial x_j = 0$, for $j=1,\ldots,n$ if there is no primal variable explicitly appears in the equation.

Dual Objective: Express primal variables in terms of $\boldsymbol{y}$ from the rest of above equation and substitute them in the Lagrange Function that becomes only in $\boldsymbol{y}$.

Warning: this may be difficult to do in general!

2-Order KKT Condition for GCO

Tangent Plane:

$T = \{\, \boldsymbol{z}: \nabla c_i(\mathbf{x})\boldsymbol{z} = 0$, for all i, such that $c_i(\boldsymbol{x})=0 \,\}$

Necessary Condition:

$\boldsymbol{z}^{\mathrm{T}}\nabla_x^{\,2}L(\mathbf{x},\mathbf{y})\boldsymbol{z} \geq \mathbf{0}$, for all $\boldsymbol{z}$ in T

Sufficient Condition:

$\boldsymbol{z}^{\mathrm{T}}\nabla_x^{\,2}L(\mathbf{x},\mathbf{y})\boldsymbol{z} > \mathbf{0}$, for all non-zero $\boldsymbol{z}$ in T

This can be done by checking positive semi-definiteness (or definiteness) of the projected Hessian of the Lagrange function

Example: Optimality Conditions

$$\begin{aligned} \min \quad & {x_1}^2 + {x_2}^2 \\ \text{s.t.} \quad & 1 - 0.25 \cdot (x_1 - 2)^2 - (x_2 - 2)^2 \geq 0 \end{aligned}$$

$$L(x_1, x_2, \lambda) = {x_1}^2 + {x_2}^2 - \lambda(1 - 0.25 \cdot (x_1 - 2)^2 - (x_2 - 2)^2)$$

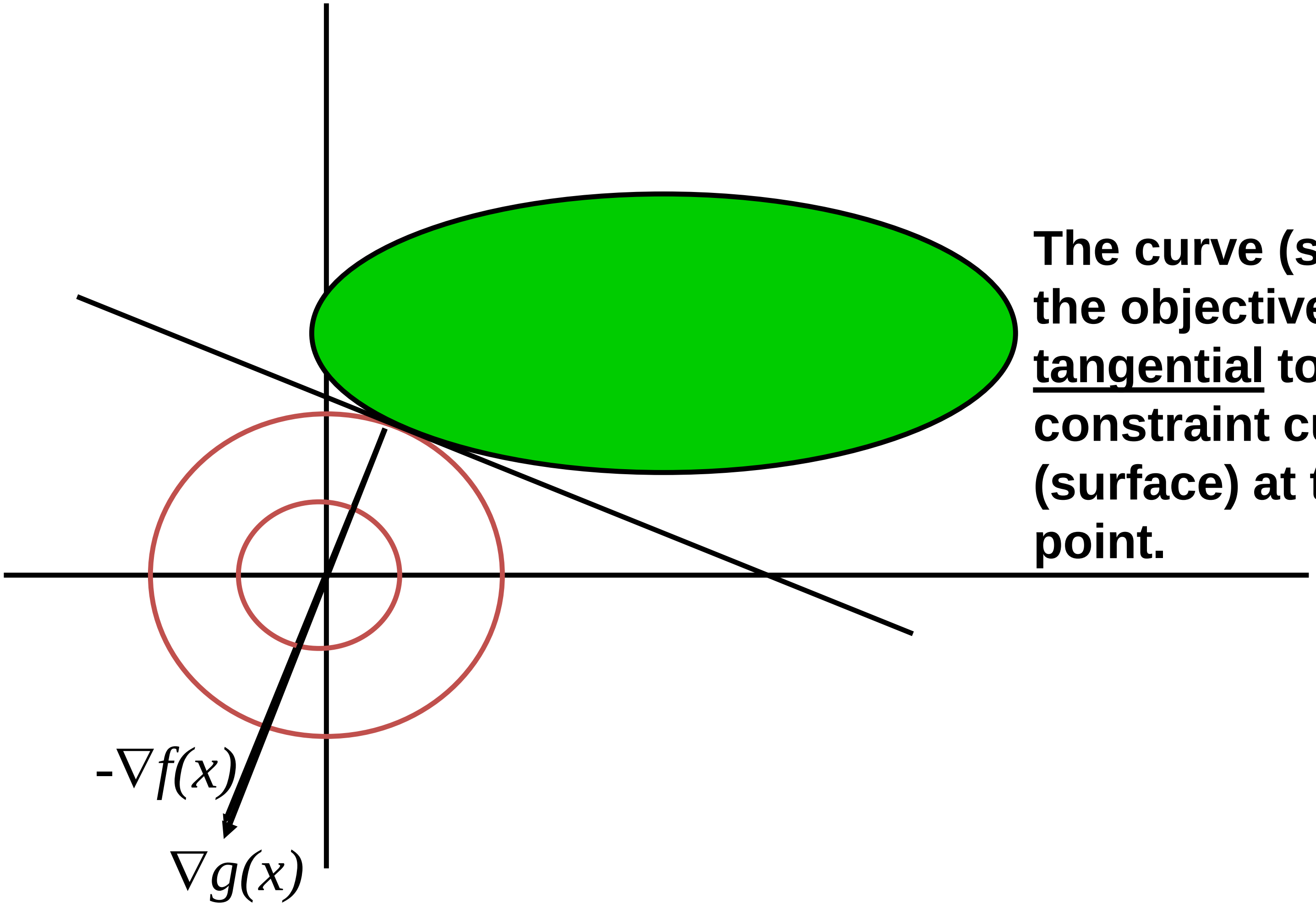
$$\begin{pmatrix} \partial L/\partial x_1 \\ \partial L/\partial x_2 \end{pmatrix} = \begin{pmatrix} 2x_1 \\ 2x_2 \end{pmatrix} - \lambda \cdot \begin{pmatrix} 0.5(2 - x_1) \\ 2(2 - x_2) \end{pmatrix} = 0, \quad \text{(LDC)}$$

$$1 - 0.25 \cdot (x_1 - 2)^2 - (x_2 - 2)^2 \geq 0, \quad \text{(OPC)}$$

$$\lambda \geq 0, \quad \text{(MSC)}$$

$$\lambda(\, 1 - 0.25 \cdot (x_1 - 2)^2 - (x_2 - 2)^2 \,) = 0. \quad \text{(CSC)}$$

Example: KKT Conditions



The curve (surface) of the objective function is <u>tangential</u> to the constraint curve (surface) at the optimal point.

Example: Computation of a KKT Point

$$\begin{pmatrix} 2x_1 \\ 2x_2 \end{pmatrix} - \lambda \cdot \begin{pmatrix} 0.5(2-x_1) \\ 2(2-x_2) \end{pmatrix} = 0$$

$$x_1 = \frac{2\lambda}{4+\lambda}; x_2 = \frac{2\lambda}{1+\lambda}$$

- If λ = 0, then x_1 = 0 and x_2 = 0, and thus the constraint would not hold with equality. Therefore, λ must be positive.
- Plugging the two values of $x_1(\lambda)$ and $x_2(\lambda)$ into the constraint with equality gives us λ = 1.8.
- We can then solve for x_1 = .61 and x_2 = 1.28.

Applications: Duality and Optimality Conditions

- The Market Equilibrium Theory
  - Zero-sum Matrix Game
  - Fisher market, Arrow-Debreu market
  - Duality and optimality lead to equilibrium conditions/computations
- Sensor localization
  - Intrinsically non-convex problems
  - SOCP relaxation: KKT conditions explain observations
  - SDP Relaxation: Duality explains localizability
- Offline and Online Convex Optimization
  - Logistic Regression and SVM
  - MDP and Reinforcement Learning
  - Learning optimal dual/Lagrange solution helps to make primal decisions online
- Robust Min-Max Optimization
  - Replace the inner max problem by its dual so that the problem becomes min-min that can be solved by any optimization solver

Sample Proofs

- Prove that, for Convex Optimization, any first-order KKT solution is a (global) minimizer (hint: use the classical Farkas' lemma).
- Prove that the singular-value computation of a given matrix is SDP-computable (hint: use the SDP solution rank theorem).
- Prove that the minimizer of a strictly convex function is unique for convex optimization.
- For a given CLO pairs where both are feasible, prove that the primal-optimal solution set is bounded if the dual has an interior feasible solution.