

Optimality Conditions for Nonlinear Optimization

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Chapters 7.1-7.3, 11.1-11.8

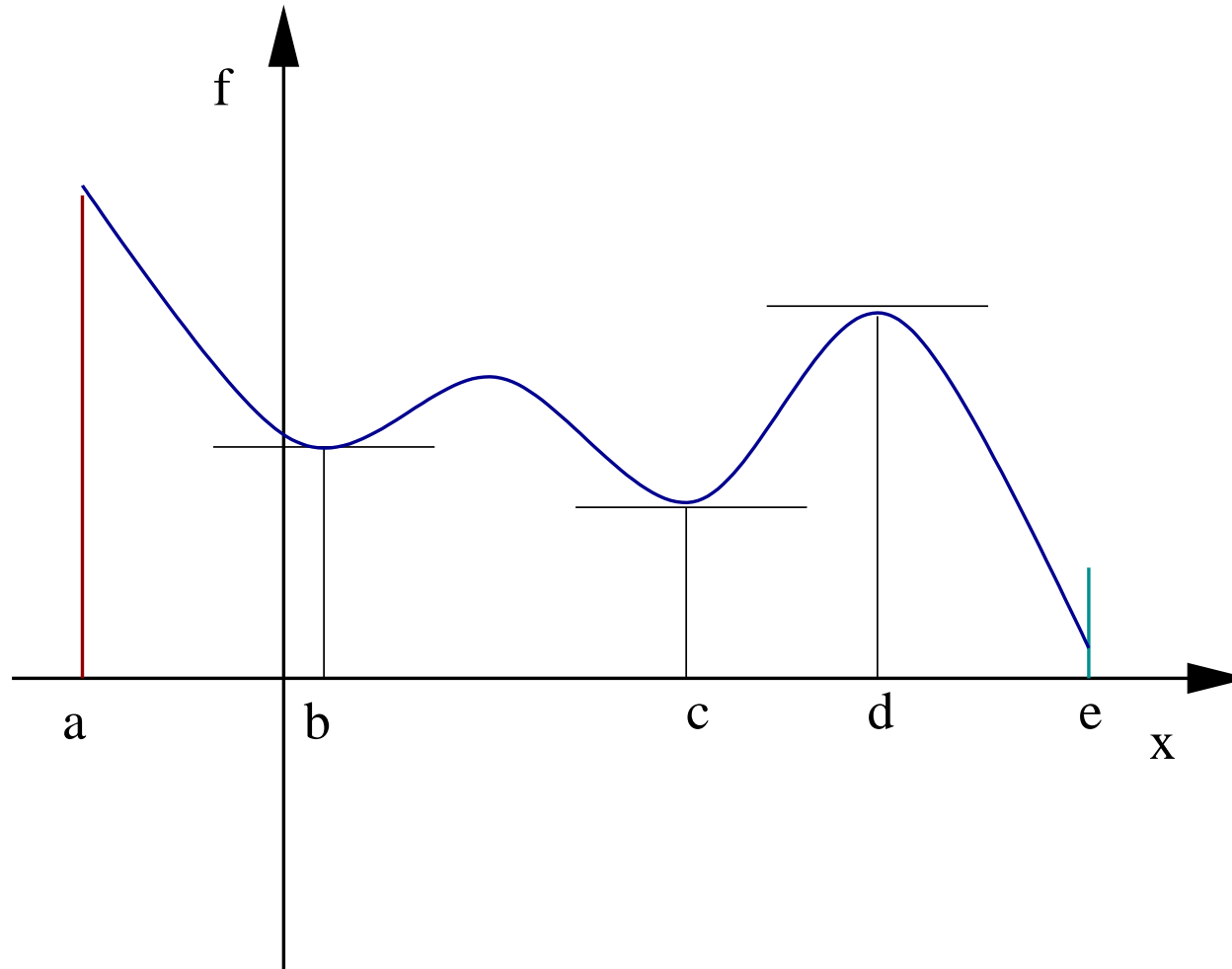
KKT Optimality Condition Illustration in One-Dimension

Figure 1: Global and Local Minimizers of One-Variable Function in Interval $[a \ e]$

A differentiable function f of one variable defined on an interval $F = [a \ e]$. If an interior-point \bar{x} is a local/global minimizer, then $f'(\bar{x}) = 0$; if the left-end-point $\bar{x} = a$ is a local minimizer, then $f'(a) \geq 0$; if the right-end-point $\bar{x} = e$ is a local minimizer, then $f'(e) \leq 0$. **first-order necessary condition (FONC)** summarizes the three cases by a unified set of **optimality/complementarity slackness conditions**:

$$a \leq x \leq e, f'(x) = y^a + y^e, y^a \geq 0, y^e \leq 0, y^a(x - a) = 0, y^e(x - e) = 0.$$

If $f'(\bar{x}) = 0$, then it is also necessary that $f(x)$ is locally convex at \bar{x} for it being a local minimizer.

How to tell the function is locally convex at the solution? It is necessary $f''(\bar{x}) \geq 0$, which is called the **second-order necessary condition (SONC)**, which we would explore further.

These conditions are still not, in general, sufficient. It does not distinguish between local minimizers, local maximizers, or saddle points.

If the **second-order sufficient condition (SOSC)**: $f''(\bar{x}) > 0$, is satisfied or the function is strictly locally convex, then \bar{x} is a local minimizer.

Thus, if the function is **convex** everywhere, the first-order necessary condition is already **sufficient**.

Second-Order Optimality Condition for Unconstrained Optimization

Theorem 1 (*First-Order Necessary Condition*) Let $f(\mathbf{x})$ be a C^1 function where $\mathbf{x} \in \mathbb{R}^n$. Then, if $\bar{\mathbf{x}}$ is a minimizer, it is necessarily $\nabla f(\bar{\mathbf{x}}) = \mathbf{0}$.

Theorem 2 (*Second-Order Necessary Condition*) Let $f(\mathbf{x})$ be a C^2 function where $\mathbf{x} \in \mathbb{R}^n$. Then, if $\bar{\mathbf{x}}$ is a minimizer, it is necessarily

$$\nabla f(\bar{\mathbf{x}}) = \mathbf{0} \quad \text{and} \quad \nabla^2 f(\bar{\mathbf{x}}) \succeq \mathbf{0}.$$

Furthermore, if $\nabla^2 f(\bar{\mathbf{x}}) \succ \mathbf{0}$, then the condition becomes *sufficient*.

The proofs would be based on 2nd-order Taylor's expansion at $\bar{\mathbf{x}}$ such that if these conditions are not satisfied, then one would be find a **descent-direction** \mathbf{d} and a small constant $\bar{\alpha} > 0$ such that $f(\bar{\mathbf{x}} + \alpha\mathbf{d}) < f(\bar{\mathbf{x}})$, $\forall 0 < \alpha \leq \bar{\alpha}$.

For example, if $\nabla f(\bar{\mathbf{x}}) = \mathbf{0}$ and $\nabla^2 f(\bar{\mathbf{x}}) \not\succeq \mathbf{0}$, the eigenvector of a negative eigenvalue of the Hessian would be a descent direction from $\bar{\mathbf{x}}$.

Again, they may still **not be sufficient**, e.g., $f(x) = x^3$.

General Optimization Problems

Let the problem have the general mathematical programming (MP) form

$$(P) \quad \begin{array}{ll} \min & f(\mathbf{x}) \\ \text{s.t.} & \mathbf{x} \in \mathcal{F}. \end{array}$$

In all forms of mathematical programming, a **feasible solution** of a given problem is a vector that satisfies the constraints of the problem, that is, in \mathcal{F} .

First question: How does one recognize or certify an optimal solution to a **generally constrained and objectived** optimization problem?

Answer: **Optimality Condition Theory**.

Descent Direction

Let f be a differentiable function on R^n . If point $\bar{\mathbf{x}} \in R^n$ and there exists a vector \mathbf{d} such that

$$\nabla f(\bar{\mathbf{x}})\mathbf{d} < 0,$$

then there exists a scalar $\bar{\tau} > 0$ such that

$$f(\bar{\mathbf{x}} + \tau\mathbf{d}) < f(\bar{\mathbf{x}}) \text{ for all } \tau \in (0, \bar{\tau}).$$

The vector \mathbf{d} (above) is called a **descent direction** at $\bar{\mathbf{x}}$. If $\nabla f(\bar{\mathbf{x}}) \neq 0$, then $\nabla f(\bar{\mathbf{x}})$ is the direction of **steepest ascent** and $-\nabla f(\bar{\mathbf{x}})$ is the direction of **steepest descent** at $\bar{\mathbf{x}}$.

Denote by $\mathcal{D}_{\bar{\mathbf{x}}}^d$ the set of descent directions at $\bar{\mathbf{x}}$, that is,

$$\mathcal{D}_{\bar{\mathbf{x}}}^d = \{\mathbf{d} \in R^n : \nabla f(\bar{\mathbf{x}})\mathbf{d} < 0\}.$$

Feasible Direction

At feasible point $\bar{\mathbf{x}}$, a feasible direction is

$$\mathcal{D}_{\bar{\mathbf{x}}}^f := \{\mathbf{d} \in \mathbb{R}^n : \mathbf{d} \neq \mathbf{0}, \bar{\mathbf{x}} + \lambda \mathbf{d} \in \mathcal{F} \text{ for all small } \lambda > 0\}.$$

Linear Constraint Examples:

$$\mathcal{F} = \mathbb{R}^n \Rightarrow \mathcal{D}^f = \mathbb{R}^n.$$

$$\mathcal{F} = \{\mathbf{x} : \mathbf{A}\mathbf{x} = \mathbf{b}\} \Rightarrow \mathcal{D}^f = \{\mathbf{d} : \mathbf{A}\mathbf{d} = \mathbf{0}\}.$$

$$\mathcal{F} = \{\mathbf{x} : \mathbf{A}\mathbf{x} \geq \mathbf{b}\} \Rightarrow \mathcal{D}^f = \{\mathbf{d} : A_i \mathbf{d} \geq 0, \forall i \in \mathcal{A}(\bar{\mathbf{x}})\},$$

where the **active** or **binding** constraint set $\mathcal{A}(\bar{\mathbf{x}}) := \{i : A_i \bar{\mathbf{x}} = b_i\}$.

Optimality Conditions

Optimality Conditions: given a feasible solution or point $\bar{\mathbf{x}}$, what are the **necessary conditions** for $\bar{\mathbf{x}}$ to be a local optimizer?

Are these conditions **sufficient**?

A general answer would be: there exists no direction at $\bar{\mathbf{x}}$ that is both **descent and feasible**. Or the **intersection** of $\mathcal{D}_{\bar{\mathbf{x}}}^d$ and $\mathcal{D}_{\bar{\mathbf{x}}}^f$ must be **empty**.

In what follows, we consider optimality conditions for Linearly Constrained Optimization Problems (LCOP).

Unconstrained Problems

Consider the **unconstrained** problem, where f is differentiable on R^n ,

$$\begin{aligned} \text{(UP)} \quad & \min && f(\mathbf{x}) \\ & \text{s.t.} && \mathbf{x} \in R^n. \end{aligned}$$

$\mathcal{D}_{\bar{\mathbf{x}}}^f = R^n$, so that $\mathcal{D}_{\bar{\mathbf{x}}}^d = \{\mathbf{d} \in R^n : \nabla f(\bar{\mathbf{x}})\mathbf{d} < 0\} = \emptyset$:

Theorem 3 Let $\bar{\mathbf{x}}$ be a (local) minimizer of (UP). If the function f is continuously differentiable at $\bar{\mathbf{x}}$, then

$$\nabla f(\bar{\mathbf{x}}) = \mathbf{0}.$$

This condition is also sufficient for global optimality if f is a convex function.

Linear Equality-Constrained Problems

Consider the **linear equality-constrained** problem, where f is differentiable on R^n ,

$$\begin{array}{ll} \text{(LEP)} & \min \quad f(\mathbf{x}) \\ & \text{s.t.} \quad A\mathbf{x} = \mathbf{b}. \end{array}$$

Theorem 4 (the Lagrange Theorem) Let $\bar{\mathbf{x}}$ be a (local) minimizer of (LEP). If the functions f is continuously differentiable at $\bar{\mathbf{x}}$, then

$$\nabla f(\bar{\mathbf{x}}) = \bar{\mathbf{y}}^T A$$

for some $\bar{\mathbf{y}} = (\bar{y}_1; \dots; \bar{y}_m) \in R^m$, which are called **Lagrange or dual multipliers**. This condition is also sufficient for global optimality if f is a convex function.

The geometric interpretation: the objective gradient vector is **perpendicular** to or the objective level set **tangents** the constraint hyperplanes.

Let $v(\mathbf{b})$ be the minimal value function of \mathbf{b} of (LEP). Then $\nabla v(\mathbf{b}) = \bar{\mathbf{y}}$.

Proof

Consider feasible direction space

$$\mathcal{F} = \{\mathbf{x} : A\mathbf{x} = \mathbf{b}\} \Rightarrow \mathcal{D}_{\bar{\mathbf{x}}}^f = \{\mathbf{d} : A\mathbf{d} = 0\}.$$

If $\bar{\mathbf{x}}$ is a local optimizer, then the **intersection** of the **descent and feasible** direction sets at $\bar{\mathbf{x}}$ must be empty or

$$A\mathbf{d} = \mathbf{0}, \nabla f(\bar{\mathbf{x}})\mathbf{d} \neq 0$$

has no feasible solution for \mathbf{d} . By **the Alternative System Theorem** it must be true that its alternative system has a solution, that is, there is $\bar{\mathbf{y}} \in R^n$ such that

$$\nabla f(\bar{\mathbf{x}}) = \bar{\mathbf{y}}^T A = \sum_{i=1}^m \bar{y}_i A_i.$$

Example: The Objective Contour Tangential to the Constraint Hyperplane

Consider the problem

$$\text{minimize} \quad (x_1 - 1)^2 + (x_2 - 1)^2$$

$$\text{subject to} \quad x_1 + x_2 = 1.$$

$$\bar{\mathbf{x}} = \begin{pmatrix} \frac{1}{2} \\ \frac{1}{2} \end{pmatrix}.$$

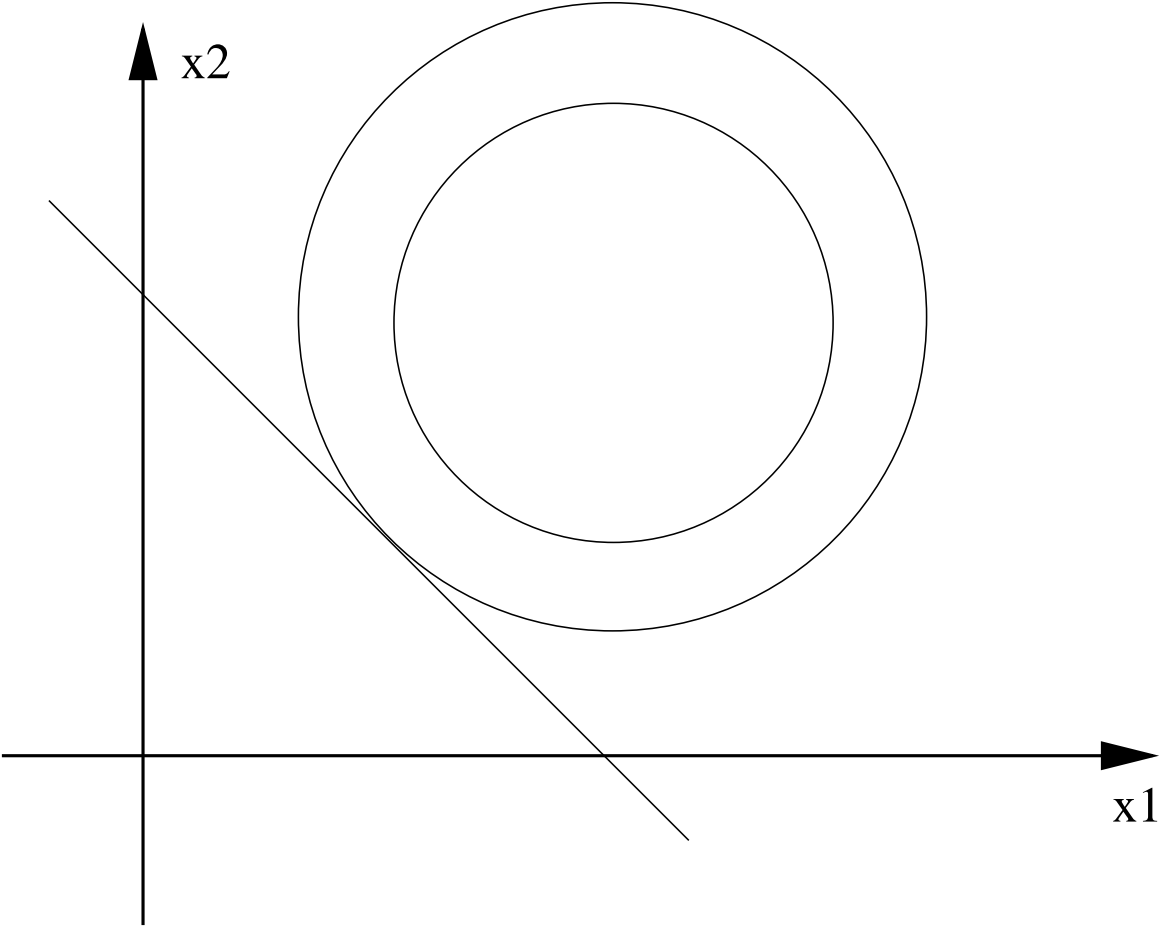


Figure 2: The Objective Contour Tangents the Constraint Hyperplane

The Barrier Optimization

Consider the problem

$$\begin{aligned} \min \quad & -\sum_{j=1}^n \log x_j \\ \text{s.t.} \quad & A\mathbf{x} = \mathbf{b}, \\ & \mathbf{x} \geq \mathbf{0} \end{aligned}$$

The non-negativity constraint can be removed if the feasible region has an "interior", that is, there is a feasible solution such that $\mathbf{x} > \mathbf{0}$. Thus, if a minimizer $\bar{\mathbf{x}}$ exists, then $\bar{\mathbf{x}} > \mathbf{0}$ and

$$-\mathbf{e}^T \bar{X}^{-1} = \bar{\mathbf{y}}^T A = \sum_{i=1}^m \bar{y}_i A_i.$$

Linear Inequality-Constrained Problems

Let us now consider the inequality-constrained problem

$$\begin{array}{ll}
 \text{(LIP)} & \min \quad f(\mathbf{x}) \\
 & \text{s.t.} \quad A\mathbf{x} \geq \mathbf{b}.
 \end{array}$$

Theorem 5 (the KKT Theorem) Let $\bar{\mathbf{x}}$ be a (local) minimizer of (LIP). If the functions f is continuously differentiable at $\bar{\mathbf{x}}$, then

$$\nabla f(\bar{\mathbf{x}}) = \bar{\mathbf{y}}^T A, \quad \bar{\mathbf{y}} \geq \mathbf{0}$$

for some $\bar{\mathbf{y}} = (\bar{y}_1; \dots; \bar{y}_m) \in R^m$, which are called *Lagrange or dual multipliers*, and $\bar{y}_i = 0$, if $i \notin \mathcal{A}(\bar{\mathbf{x}})$. These conditions are also sufficient for global optimality if f is a convex function.

The geometric interpretation: the objective gradient vector is in the **cone** generated by the **normal directions** of the active-constraint hyperplanes.

Proof

$$\mathcal{F} = \{\mathbf{x} : \mathbf{Ax} \geq \mathbf{b}\} \Rightarrow \mathcal{D}_{\bar{\mathbf{x}}}^f = \{\mathbf{d} : A_i \mathbf{d} \geq 0, \forall i \in \mathcal{A}(\bar{\mathbf{x}})\},$$

or

$$\mathcal{D}_{\bar{\mathbf{x}}}^f = \{\mathbf{d} : \bar{\mathbf{A}}\mathbf{d} \geq \mathbf{0}\},$$

where $\bar{\mathbf{A}}$ corresponds to those active constraints. If $\bar{\mathbf{x}}$ is a local optimizer, then the **intersection** of the **descent and feasible** direction sets at $\bar{\mathbf{x}}$ must be empty or

$$\bar{\mathbf{A}}\mathbf{d} \geq \mathbf{0}, \nabla f(\bar{\mathbf{x}})\mathbf{d} < 0$$

has no feasible solution. By **the Alternative System Theorem** it must be true that its alternative system has a solution, that is, there is $\bar{\mathbf{y}} \geq \mathbf{0}$ such that

$$\nabla f(\bar{\mathbf{x}}) = \bar{\mathbf{y}}^T \bar{\mathbf{A}} = \sum_{i \in \mathcal{A}(\bar{\mathbf{x}})} \bar{y}_i A_i = \sum_i \bar{y}_i A_i,$$

when let $\bar{y}_i = 0$ for all $i \notin \mathcal{A}(\bar{\mathbf{x}})$. Then we prove the theorem.

Example: The Gradient is in the Normal Cone of the Half Spaces

Consider the problem

$$\min (x_1 - 1)^2 + (x_2 - 1)^2$$

$$\text{s.t.} \quad -x_1 - 2x_2 \geq -1,$$

$$-2x_1 - x_2 \geq -1.$$

$$\bar{\mathbf{x}} = \left(\frac{1}{3}; \frac{1}{3} \right).$$

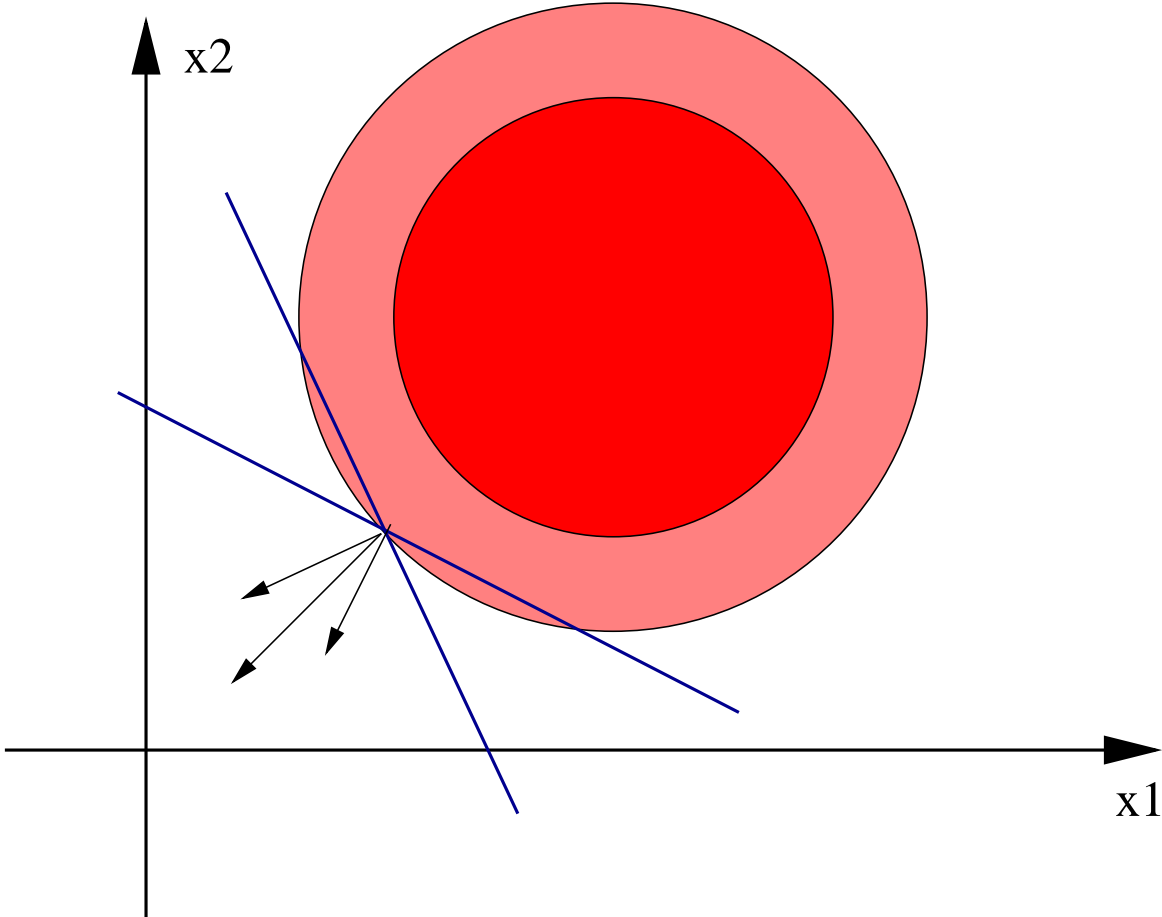


Figure 3: The objective gradient in the normal cone of the half spaces

Optimization with Mixed Linear Constraints

We now consider optimality conditions for problems having both **inequality and equality** constraints. These can be denoted

$$(P) \quad \begin{array}{ll} \min & f(\mathbf{x}) \\ \text{s.t.} & \mathbf{a}_i \mathbf{x} \quad (\leq, =, \geq) \quad b_i, \quad i = 1, \dots, m, \end{array}$$

For any feasible point $\bar{\mathbf{x}}$ of (P) we have the sets

$$\mathcal{A}(\bar{\mathbf{x}}) = \{i : \mathbf{a}_i \bar{\mathbf{x}} = b_i\}$$

$$\mathcal{D}_{\bar{\mathbf{x}}}^d = \{\mathbf{d} : \nabla f(\bar{\mathbf{x}}) \mathbf{d} < 0\}.$$

The KKT Theorem Summary

Theorem 6 Let $\bar{\mathbf{x}}$ be a local minimizer for (P). Then there exist multipliers $\bar{\mathbf{y}}$ such that

$$\begin{array}{llll}
 \mathbf{a}_i \bar{\mathbf{x}} & (\leq, =, \geq) & b_i, \quad i = 1, \dots, m, & \text{(Original Problem Constraints (OPC))} \\
 \nabla f(\bar{\mathbf{x}}) & = & \bar{\mathbf{y}}^T \mathbf{A} & \text{(Lagrangian Multiplier Conditions (LMC))} \\
 \bar{y}_i & (\leq, 'free', \geq) & 0, \quad i = 1, \dots, m, & \text{(Multiplier Sign Constraints (MSC))} \\
 \bar{y}_i & = & 0 \quad \text{if } i \notin \mathcal{A}(\bar{\mathbf{x}}) & \text{(Complementarity Slackness Conditions (CSC))}.
 \end{array}$$

These conditions are also sufficient for global optimality if f is a convex function.

When First-Order Optimality Conditions are Sufficient?

Theorem 7 If objective f is a locally *convex* function in the feasible direction space at the KKT solution $\bar{\mathbf{x}}$, then the (first-order) KKT optimality conditions are *sufficient* for the *local optimality* at $\bar{\mathbf{x}}$.

A function is locally convex in a space D means that $\phi(\alpha) := f(\bar{\mathbf{x}} + \alpha \mathbf{d})$ is a convex function of α in a sufficiently small neighborhood of 0 for all $\mathbf{d} \in D$.

Corollary 1 If f is differentiable *convex* function in the feasible region, then the (first-order) KKT optimality conditions are *sufficient* for the *global optimality* for linearly constrained optimization.

Again, how to check convexity, say $f(x) = x^3$?

- Hessian matrix is PSD in the feasible region.
- Epigraph is a convex set.

LCOP Examples: Linear Optimization

$$\begin{aligned}
 (LP) \quad & \min \quad \mathbf{c}^T \mathbf{x} \\
 & \text{s.t.} \quad A\mathbf{x} = \mathbf{b}, \mathbf{x} \geq \mathbf{0}.
 \end{aligned}$$

For any feasible \mathbf{x} of (LP), it's optimal if for some \mathbf{y}, \mathbf{s}

$$\begin{aligned}
 x_j s_j &= 0, \quad \forall j = 1, \dots, n \\
 A\mathbf{x} &= \mathbf{b} \\
 \nabla(\mathbf{c}^T \mathbf{x}) = \mathbf{c}^T &= \mathbf{y}^T A + \mathbf{s}^T \\
 \mathbf{x}, \mathbf{s} &\geq \mathbf{0}.
 \end{aligned}$$

Here, \mathbf{y} (shadow prices in LP) are Lagrange or dual multipliers of equality constraints, and \mathbf{s} (reduced gradient/costs in LP) are Lagrange or dual multipliers for $\mathbf{x} \geq \mathbf{0}$.

LCOP Examples : Quadratic Optimization

$$\begin{aligned} (QP) \quad & \min \quad \mathbf{x}^T Q \mathbf{x} - 2\mathbf{c}^T \mathbf{x} \\ & \text{s.t.} \quad A\mathbf{x} = \mathbf{b}, \mathbf{x} \geq \mathbf{0}. \end{aligned}$$

Optimality Conditions:

$$\begin{aligned} x_j s_j &= 0, \quad \forall j = 1, \dots, n \\ A\mathbf{x} &= \mathbf{b} \\ 2Q\mathbf{x} - 2\mathbf{c} - A^T \mathbf{y} - \mathbf{s} &= \mathbf{0} \\ \mathbf{x}, \mathbf{s} &\geq \mathbf{0} \end{aligned}$$

LCOP Examples: Linear Barrier Optimization

$$\min f(\mathbf{x}) = \mathbf{c}^T \mathbf{x} - \mu \sum_{j=1}^n \log(x_j), \text{ s.t. } A\mathbf{x} = \mathbf{b}, \mathbf{x} \geq \mathbf{0}.$$

for some fixed $\mu > 0$. Assume that interior of the feasible region is not empty:

$$\begin{aligned} A\mathbf{x} &= \mathbf{b} \\ c_j - \frac{\mu}{x_j} - (\mathbf{y}^T A)_j &= 0, \forall j = 1, \dots, n \\ \mathbf{x} &> \mathbf{0}. \end{aligned}$$

Let $s_j = \frac{\mu}{x_j}$ for all j (note that this \mathbf{s} is not the \mathbf{s} in the KKT condition of $f(\mathbf{x})$). Then

$$\begin{aligned} x_j s_j &= \mu, \forall j = 1, \dots, n, \\ A\mathbf{x} &= \mathbf{b}, \\ A^T \mathbf{y} + \mathbf{s} &= \mathbf{c}, \\ (\mathbf{x}, \mathbf{s}) &> \mathbf{0}. \end{aligned}$$

KKT Application: Fisher's Equilibrium Prices

Player $i \in B$'s optimization problem for given prices $p_j, j \in G$.

$$\begin{aligned} \max \quad & \mathbf{u}_i^T \mathbf{x}_i := \sum_{j \in G} u_{ij} x_{ij} \\ \text{s.t.} \quad & \mathbf{p}^T \mathbf{x}_i := \sum_{j \in G} p_j x_{ij} \leq w_i, \\ & x_{ij} \geq 0, \quad \forall j, \end{aligned}$$

Assume that the given amount of each good is \bar{s}_j . The equilibrium price vector is the one that for all $j \in G$

$$\sum_{i \in B} x^*(\mathbf{p})_{ij} = \bar{s}_j$$

where $\mathbf{x}^*(\mathbf{p})$ is a maximizer of the utility maximization problem for every buyer i .

Example of Fisher's Equilibrium Prices

There two goods, x and y , each with 1 unit on the market. Buyer 1, 2's optimization problems for given prices p_x, p_y .

$$\begin{aligned} \max \quad & 2x_1 + y_1 \\ \text{s.t.} \quad & p_x \cdot x_1 + p_y \cdot y_1 \leq 5, \\ & x_1, y_1 \geq 0; \end{aligned}$$

$$\begin{aligned} \max \quad & 3x_2 + y_2 \\ \text{s.t.} \quad & p_x \cdot x_2 + p_y \cdot y_2 \leq 8, \\ & x_2, y_2 \geq 0. \end{aligned}$$

$$p_x = \frac{26}{3}, p_y = \frac{13}{3}, x_1 = \frac{1}{13}, y_1 = 1, x_2 = \frac{12}{13}, y_2 = 0$$

Equilibrium Price Conditions

Player $i \in B$'s dual problem for given prices $p_j, j \in G$.

$$\begin{aligned} \min \quad & w_i y_i \\ \text{s.t.} \quad & \mathbf{p} y_i \geq \mathbf{u}_i, y_i \geq 0 \end{aligned}$$

The necessary and sufficient conditions for an equilibrium point \mathbf{x}_i, \mathbf{p} are:

$$\begin{aligned} \mathbf{p}^T \mathbf{x}_i = w_i, \mathbf{x}_i \geq \mathbf{0}, \quad \forall i, \\ p_j y_i \geq u_{ij}, y_i \geq 0, \quad \forall i, j, \\ \mathbf{u}_i^T \mathbf{x}_i = w_i y_i, \quad \forall i, \\ \sum_i x_{ij} = \bar{s}_j, \quad \forall j. \end{aligned} \quad \text{or} \quad \begin{aligned} \mathbf{p}^T \mathbf{x}_i = w_i, \mathbf{x}_i \geq \mathbf{0}, \quad \forall i, \\ p_j \frac{\mathbf{u}_i^T \mathbf{x}_i}{w_i} \geq u_{ij}, \quad \forall i, j, \\ \sum_i x_{ij} = \bar{s}_j, \quad \forall j. \end{aligned}$$

Equilibrium Price Conditions continued

These conditions can be further simplified to

$$\begin{aligned}\sum_j \bar{s}_j p_j &= \sum_i w_i, \quad \mathbf{x}_i \geq \mathbf{0}, \quad \forall i, \\ p_j &\geq u_{ij} \frac{w_i}{\mathbf{u}_i^T \mathbf{x}_i}, \quad \forall i, j, \\ \sum_i x_{ij} &= \bar{s}_j, \quad \forall j.\end{aligned}$$

since from the second inequality (after multiplying x_{ij} to both sides and take sum over j) we have

$$\mathbf{p}^T \mathbf{x}_i \geq w_i, \quad \forall i.$$

Then, from the rest conditions

$$\sum_i w_i = \sum_j \bar{s}_j p_j = \sum_i \mathbf{p}^T \mathbf{x}_i \geq \sum_i w_i.$$

Thus, these conditions imply $\mathbf{p}^T \mathbf{x}_i = w_i, \quad \forall i.$

Equilibrium Price Property

If u_{ij} has at least one positive coefficient for every j , then we must have $p_j > 0$ for every j at every equilibrium. Moreover, The second inequality can be rewritten as

$$\log(\mathbf{u}_i^T \mathbf{x}_i) + \log(p_j) \geq \log(w_i) + \log(u_{ij}), \quad \forall i, j, u_{ij} > 0.$$

The function on the left is (strictly) concave in \mathbf{x}_i and p_j . Thus,

Theorem 8 *The equilibrium set of the Fisher Market is convex.*

Aggregated Social Optimization

$$\begin{aligned} \max \quad & \sum_{i \in B} w_i \log(\mathbf{u}_i^T \mathbf{x}_i) \\ \text{s.t.} \quad & \sum_{i \in B} x_{ij} = \bar{s}_j, \quad \forall j \in G \\ & x_{ij} \geq 0, \quad \forall i, j, \end{aligned}$$

Theorem 9 (Eisenberg and Gale 1959) *Optimal dual (Lagrange) multiplier vector of equality constraints is an equilibrium price vector.*

Proof: The optimality conditions of the social problem are **identical** to the equilibrium conditions.

Aggregated Example

$$\begin{aligned} \max \quad & 5 \log(2x_1 + y_1) + 8 \log(3x_2 + y_2) \\ \text{s.t.} \quad & x_1 + x_2 = 1, \\ & y_1 + y_2 = 1, \\ & x_1, x_2, y_1, y_2 \geq 0. \end{aligned}$$

Or

$$\begin{aligned} \max \quad & 5 \log(u_1) + 8 \log(u_2) \\ \text{s.t.} \quad & 2x_1 + y_1 - u_1 = 0, \\ & 3x_2 + y_2 - u_2 = 0, \\ & x_1 + x_2 = 1, \\ & y_1 + y_2 = 1, \\ & x_1, x_2, y_1, y_2 \geq 0. \end{aligned}$$

Optimality Conditions of the Aggregated Problem

$$\begin{aligned}
 w_i \frac{u_{ij}}{\mathbf{u}_i^T \mathbf{x}_i} &\leq p_j, \quad \forall i, j \\
 w_i \frac{u_{ij} x_{ij}}{\mathbf{u}_i^T \mathbf{x}_i} &= p_j x_{ij}, \quad \forall i, j \quad (\text{complementarity}) \\
 \sum_i x_{ij} &= \bar{s}_j, \quad \forall j \\
 \mathbf{x}_i, \mathbf{p} &\geq \mathbf{0}.
 \end{aligned}$$

Let $y_i = \frac{\mathbf{u}_i^T \mathbf{x}_i}{w_i}$. Then, these conditions are **identical** to the equilibrium price conditions, since

$$y_i = \frac{\mathbf{u}_i^T \mathbf{x}_i}{w_i} \geq \frac{u_{ij}}{p_j}, \quad \forall i, j.$$

Question: Is the price vector \mathbf{p} **unique** when at least one $u_{ij} > 0$ among $i \in B$ and $u_{ij} > 0$ among $j \in G$.

KKT Application: “Inverse” Learning

Suppose the market does not know the “private information” of each buyer, but can observe realized activities x_{ij}^t from each buyer and the market prices \mathbf{p}^t from various known supply date set $\bar{\mathbf{s}}^t$, $t = 1, \dots, T$. Could we learn the private information u_{ij} (partially) and w_i^t from each buyer?

How should we formulate the learning problem?

Using Deep Learning?

Using KKT/Equilibrium conditions?