

School of Engineering

# CG and MINRES: An empirical comparison

## Prequel to LSQR and LSMR: Two least-squares solvers

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# Abstract

For iterative solution of symmetric systems  $Ax = b$ , the conjugate gradient method (CG) is commonly used when  $A$  is positive definite, while the minimal residual method (MINRES) is typically reserved for indefinite systems. We investigate the sequence of solutions generated by each method and suggest that **even if  $A$  is positive definite, MINRES may be preferable to CG** if iterations are to be terminated early.

The classic symmetric positive-definite system comes from the full-rank least-squares (LS) problem  $\min \|Ax - b\|$ . Specialization of CG and MINRES to the associated normal equation  $A^T Ax = A^T b$  leads to LSQR and LSMR respectively. We include numerical comparisons of these two LS solvers because they motivated this retrospective study of CG versus MINRES.

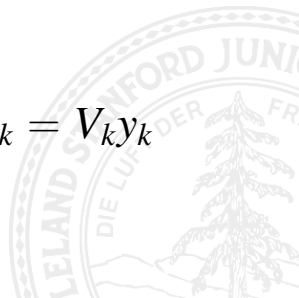
- ① CG and MINRES
  - The Lanczos Process
  - Properties
  - Backward Errors
- ② LSQR and LSMR
- ③ LSMR Derivation
  - Golub-Kahan bidiagonalization
  - Properties
- ④ LSMR Experiments
  - Backward Errors
- ⑤ Summary



# Part I: CG and MINRES

Iterative algorithms for  $Ax = b$ ,  $A = A^T$   
based on the Lanczos process

Krylov-subspace methods:  $x_k = V_k y_k$



## Lanczos process (summary)

$$\beta_1 v_1 = b \quad AV_k = V_{k+1} H_k$$

$$V_k = (v_1 \ v_2 \ \cdots \ v_k)$$

$$T_k = \begin{pmatrix} \alpha_1 & \beta_2 & & & \\ \beta_2 & \alpha_2 & \ddots & & \\ & \ddots & \ddots & \beta_k & \\ & & & \beta_k & \alpha_k \end{pmatrix}$$

$$H_k = \begin{pmatrix} T_k \\ \beta_{k+1} e_k^T \end{pmatrix}$$



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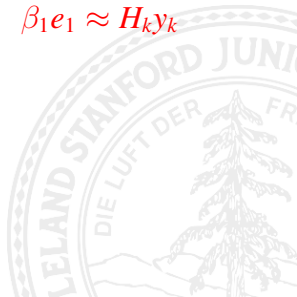
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$$H_k = \begin{pmatrix} T_k \\ \beta_{k+1} e_k^T \end{pmatrix}$$

$$\begin{aligned} r_k &= b - Ax_k \\ &= \beta_1 v_1 - AV_k y_k \\ &= V_{k+1} (\beta_1 e_1 - H_k y_k), \end{aligned}$$

Aim:  $\beta_1 e_1 \approx H_k y_k$



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$$= V_{k+1} (\beta_1 e_1 - H_k y_k),$$

Aim:  $\beta_1 e_1 \approx H_k y_k$

## Two subproblems

CG

$$T_k y_k = \beta_1 e_1$$

$$x_k = V_k y_k$$

MINRES

$$\min \|H_k y_k - \beta_1 e_1\|$$

$$x_k = V_k y_k$$

## Common practice

$$Ax = b, \quad A = A^T$$





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$A$  positive definite  $\Rightarrow$  Use CG

$A$  indefinite  $\Rightarrow$  Use MINRES



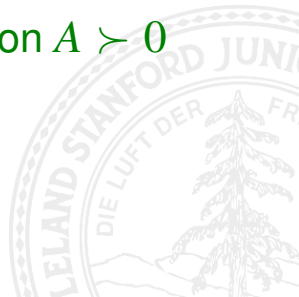
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Experiment: CG vs MINRES on  $A \succ 0$



## Common practice

$$Ax = b, \quad A = A^T$$

$A$  positive definite  $\Rightarrow$  Use CG

$A$  indefinite  $\Rightarrow$  Use MINRES

### Experiment: CG vs MINRES on $A \succ 0$

- Hestenes and Stiefel (1952) proposed both CG and CR for  $A \succ 0$  and proved many properties
- $CR \equiv MINRES$  when  $A \succ 0$   
They both minimize  $\|r_k\| = \|b - Ax_k\|$  in the Krylov subspace

# Theoretical properties for $Ax = b, A \succ 0$

		CG	CR (MINRES)
$\ x^* - x_k\ $	↘	HS 1952	HS 1952
$\ x^* - x_k\ _A$	↘	HS 1952	HS 1952
$\ x_k\ $	↗	Steihaug 1983	Fong 2012



# Theoretical properties for $Ax = b, A \succ 0$

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		CR (MINRES)
$\ r_k\ $	↘	HS 1952
$\ r_k\  / \ x_k\ $	↘	Fong 2012

# Backward error for square systems $Ax = b$

An approximate solution  $x_k$  is **acceptable** iff  $\exists E, f$  st

$$(A + E)x_k = b + f \quad \frac{\|E\|}{\|A\|} \leq \alpha \quad \frac{\|f\|}{\|b\|} \leq \beta$$



# Backward error for square systems $Ax = b$

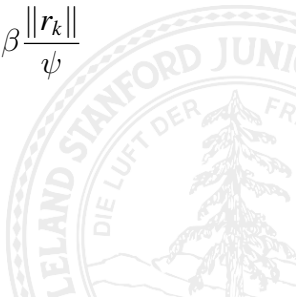
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Smallest perturbations  $E, f$ : (Tittley-Peloquin 2010)

$$E = \frac{\alpha \|A\|}{\psi \|x_k\|} r_k x_k^T \quad \frac{\|E\|}{\|A\|} = \alpha \frac{\|r_k\|}{\psi}$$

$$f = -\frac{\beta \|b\|}{\psi} r_k \quad \frac{\|f\|}{\|b\|} = \beta \frac{\|r_k\|}{\psi}$$



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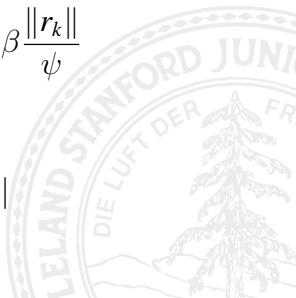
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$$f = -\frac{\beta \|b\|}{\psi} r_k \quad \frac{\|f\|}{\|b\|} = \beta \frac{\|r_k\|}{\psi}$$

Stopping rule:

$$\|r_k\| \leq \psi \equiv \alpha \|A\| \|x_k\| + \beta \|b\|$$





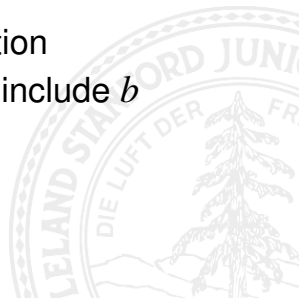
# Backward error for square systems, $\beta = 0$

$$(A + E^{(k)})x_k = b$$
$$E^{(k)} = \frac{r_k x_k^T}{\|x_k\|^2} \quad \|E^{(k)}\| = \frac{\|r_k\|}{\|x_k\|}$$

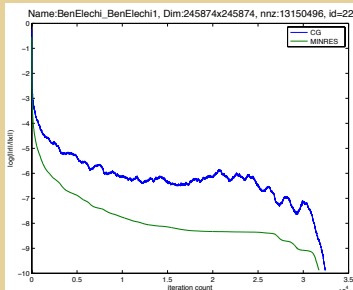
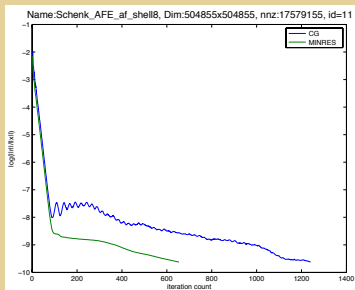
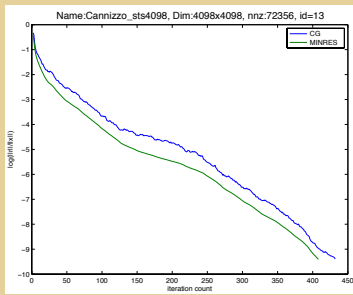
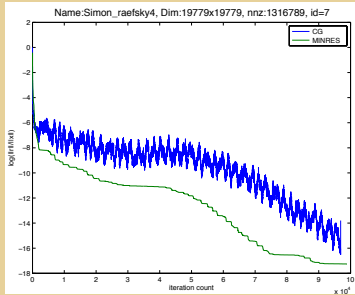
Data: Tim Davis's sparse matrix collection

Real, symmetric posdef examples that include  $b$

Plot  $\log_{10} \|E^{(k)}\|$  for CG and MINRES



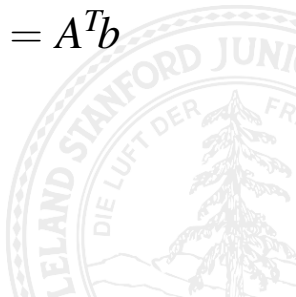
# Backward Error of CG vs MINRES on $A \succ 0$



## Part II: LSQR and LSMR

LSQR  $\equiv$  CG      on  $A^T A x = A^T b$

LSMR  $\equiv$  MINRES on  $A^T A x = A^T b$



# What problems do LSQR and LSMR solve?

solve  $Ax = b$



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$$\begin{array}{ll} \min & \|x\| \\ \text{st} & Ax = b \end{array}$$

$$\min \left\| \begin{pmatrix} A \\ \lambda I \end{pmatrix} x - \begin{pmatrix} b \\ 0 \end{pmatrix} \right\|_2$$



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$$\min \|Ax - b\|_2$$

$$\begin{array}{ll} \min & \|x\| \\ \text{st} & Ax = b \end{array}$$

$$\min \left\| \begin{pmatrix} A \\ \lambda I \end{pmatrix} x - \begin{pmatrix} b \\ 0 \end{pmatrix} \right\|_2$$

## Properties

- $A$  is rectangular ( $m \times n$ ) and often sparse
- $A$  can be an operator ( $\Rightarrow$  allows preconditioning)
- $Av, A^T u$  plus  $O(m + n)$  operations per iteration



# Why invent another algorithm?



Reason one

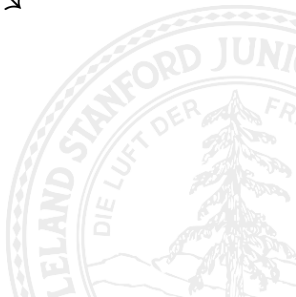
# CG vs MINRES



Reason two

# Monotone convergence of residuals

$$\|r_k\| \quad \text{and} \quad \|A^T r_k\| \quad \searrow$$



$$\min \|Ax - b\|$$

Measure of convergence

- $r_k = b - Ax_k$
- $\|r_k\| \rightarrow \|\hat{r}\|, \|A^T r_k\| \rightarrow 0$

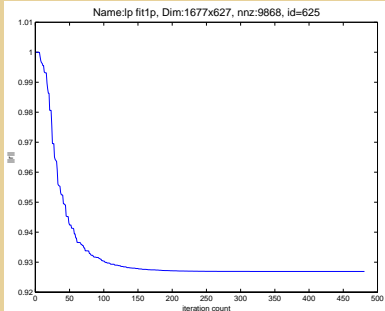


$$\min \|Ax - b\|$$

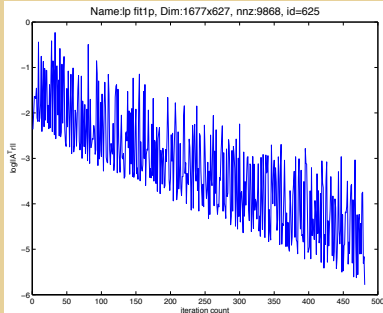
Measure of convergence

- $r_k = b - Ax_k$
- $\|r_k\| \rightarrow \|\hat{r}\|, \|A^T r_k\| \rightarrow 0$

LSQR  $\|r_k\|$



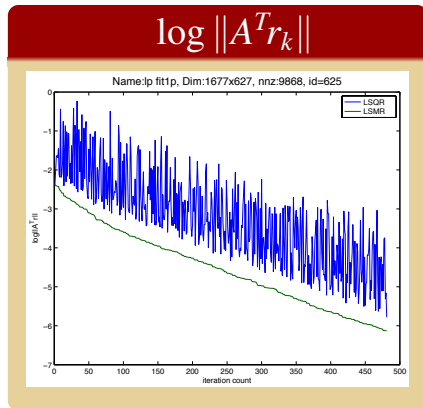
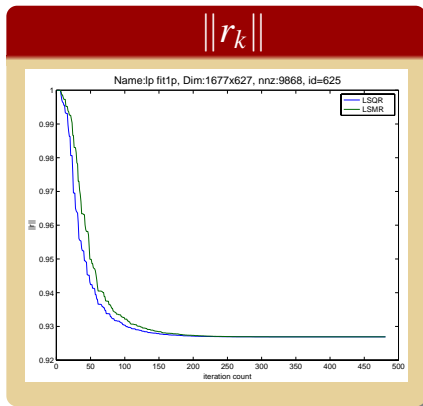
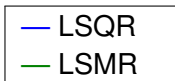
LSQR  $\log \|A^T r_k\|$



$$\min \|Ax - b\|$$

Measure of convergence

- $r_k = b - Ax_k$
- $\|r_k\| \rightarrow \|\hat{r}\|, \|A^T r_k\| \rightarrow 0$



# LSMR Derivation



# Golub-Kahan bidiagonalization

Given  $A$  ( $m \times n$ ) and  $b$  ( $m \times 1$ )

## Direct bidiagonalization

$$U^T (b \ A) \begin{pmatrix} 1 \\ \mathbf{v} \end{pmatrix} = \begin{pmatrix} \times & \times & & \\ & \times & \times & \\ & & \times & \times \\ & & & \times & \times \end{pmatrix} \Rightarrow (b \ AV) = U (\beta_1 e_1 \ B)$$





# Golub-Kahan bidiagonalization

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## Iterative bidiagonalization $\text{Bidiag}(A, b)$

Half a page in the 1965 Golub-Kahan SVD paper

## Golub-Kahan bidiagonalization (2)

$$b = U_{k+1}(\beta_1 e_1)$$

$$AV_k = U_{k+1}B_k$$

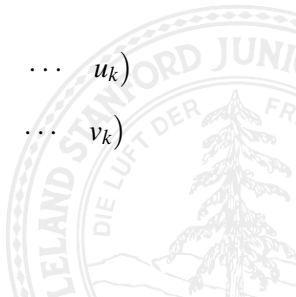
$$A^T U_k = V_k B_k^T \begin{pmatrix} I_k \\ 0 \end{pmatrix}$$

where

$$B_k = \begin{pmatrix} \alpha_1 & & & & & \\ \beta_2 & \alpha_2 & & & & \\ & \ddots & \ddots & & & \\ & & & \beta_k & \alpha_k & \\ & & & & \beta_{k+1} & \end{pmatrix}$$

$$U_k = (u_1 \quad \cdots \quad u_k)$$

$$V_k = (v_1 \quad \cdots \quad v_k)$$



## Golub-Kahan bidiagonalization (3)

$V_k$  spans the Krylov subspace:

$$\text{span}\{v_1, \dots, v_k\} = \text{span}\{A^T b, (A^T A)A^T b, \dots, (A^T A)^{k-1}A^T b\}$$



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Define  $x_k = V_k y_k$

### Subproblem to solve

$$\min_{y_k} \|r_k\| = \min_{y_k} \|\beta_1 e_1 - B_k y_k\| \quad (\text{LSQR})$$

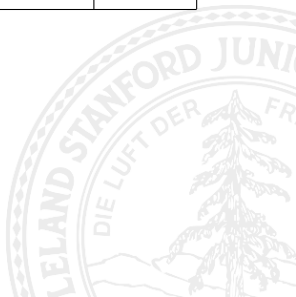
$$\min_{y_k} \|A^T r_k\| = \min_{y_k} \left\| \bar{\beta}_1 e_1 - \begin{pmatrix} B_k^T B_k \\ \bar{\beta}_{k+1} e_k^T \end{pmatrix} y_k \right\| \quad (\text{LSMR})$$

where  $r_k = b - Ax_k$ ,  $\bar{\beta}_k = \alpha_k \beta_k$

# Computational and storage requirement

	Storage		Work	
	$m$	$n$	$m$	$n$
MINRES on $A^T A x = A^T b$	$Av_1$	$x, v_1, v_2, w_1, w_2$		8
LSQR	$Av, u$	$x, v, w$	3	5
LSMR	$Av, u$	$x, v, h, \bar{h}$	3	6

where  $h_k, \bar{h}_k$  are scalar multiples of  $w_k, \bar{w}_k$



# Theoretical properties for $\min \|Ax = b\|$

		LSQR	LSMR
$\ x^* - x_k\ $	↘	HS 1952	HS 1952
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$\ x_k\ $	↗	Steihaug 1983	Fong 2012
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		LSQR	LSMR
$\ A^T r_k\ $	↘		FS 2011
$\ A^T r_k\  / \ r_k\ $	↘		<b>mostly</b>
		$\ A^T r_k\  / \ r_k\  \geq$	$\ A^T r_k\  / \ r_k\ $

# LSMR Experiments





# Overdetermined systems

## Test Data

- Tim Davis, University of Florida Sparse Matrix Collection
- LPnetlib: Linear Programming Problems
- $A = (\text{Problem.A})'$     $b = \text{Problem.c}$    (127 problems)



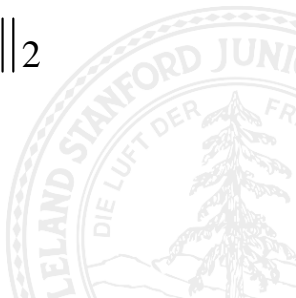
# Overdetermined systems

## Test Data

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- $A = (\text{Problem.A})'$     $b = \text{Problem.c}$    (127 problems)

$$\text{Solve } \min \|Ax - b\|_2$$

with LSQR and LSMR



# Backward error – estimates

$$\begin{array}{ll} A^T A \hat{x} = A^T b & \hat{r} = b - A \hat{x} \quad \text{exact} \\ (A + E_i)^T (A + E_i) x = (A + E_i)^T b & r = b - A x \quad \text{any } x \end{array}$$



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 \end{array}$$

Two estimates given by Stewart (1975 and 1977)

$$\begin{array}{lll}
 E_1 = \frac{ex^T}{\|x\|^2} & \|E_1\| = \frac{\|e\|}{\|x\|} & e = \hat{r} - r \\
 E_2 = -\frac{rr^T A}{\|r\|^2} & \|E_2\| = \frac{\|A^T r\|}{\|r\|} & \text{computable}
 \end{array}$$



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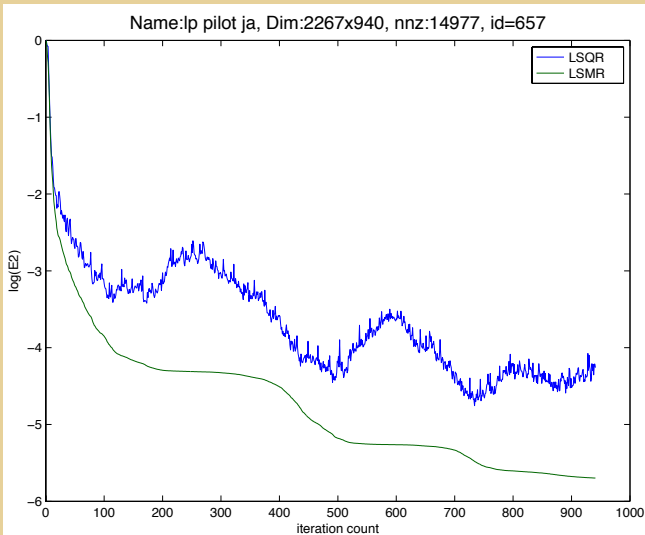
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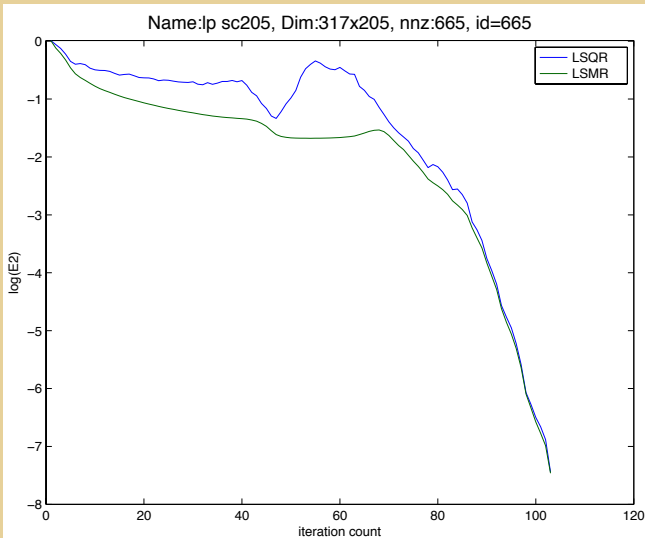
### Theorem

$$\|E_2^{\text{LSMR}}\| \leq \|E_2^{\text{LSQR}}\|$$

# $\log_{10} \|E_2\|$ for LSQR and LSMR – typical



# $\log_{10} \|E_2\|$ for LSQR and LSMR – rare



## Backward error - optimal

$$\mu(x) \equiv \min_E \|E\| \quad \text{st} \quad (A + E)^T(A + E)x = (A + E)^T b$$

Exact  $\mu(x)$  (Waldén, Karlson, & Sun 1995, Higham 2002)

$$C \equiv \left[ A \quad \frac{\|r\|}{\|x\|} \left( I - \frac{rr^T}{\|r\|^2} \right) \right] \quad \mu(x) = \sigma_{\min}(C)$$





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$$\mu(x) \equiv \min_E \|E\| \quad \text{st} \quad (A + E)^T(A + E)x = (A + E)^T b$$

Cheaper estimate  $\tilde{\mu}(x)$  (Grcar, Saunders, & Su 2007)

$$K = \begin{pmatrix} A \\ \frac{\|r\|}{\|x\|} I \end{pmatrix} \quad v = \begin{pmatrix} r \\ 0 \end{pmatrix}$$
$$\min_y \|Ky - v\| \quad \tilde{\mu}(x) = \frac{\|Ky\|}{\|x\|}$$



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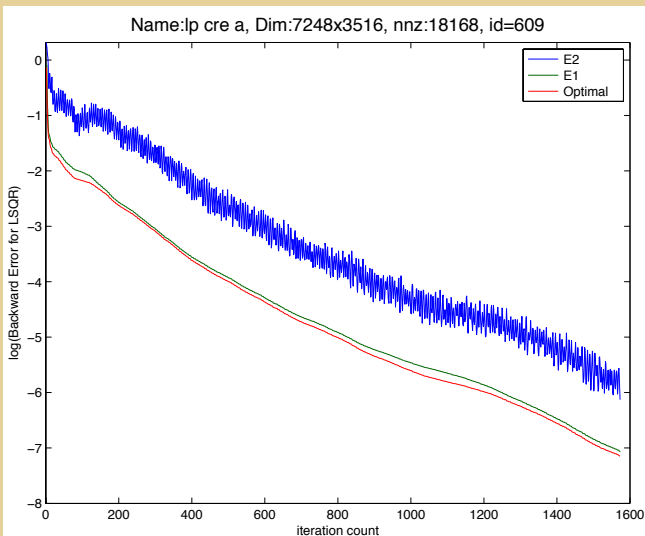
```

r = b - A*x;
p = colamd(A);
eta = norm(r)/norm(x);
K = [A(:,p); eta*speye(n)];
v = [r; zeros(n,1)];
[c,R] = qr(K,v,0);
mutilde = norm(c)/norm(x);

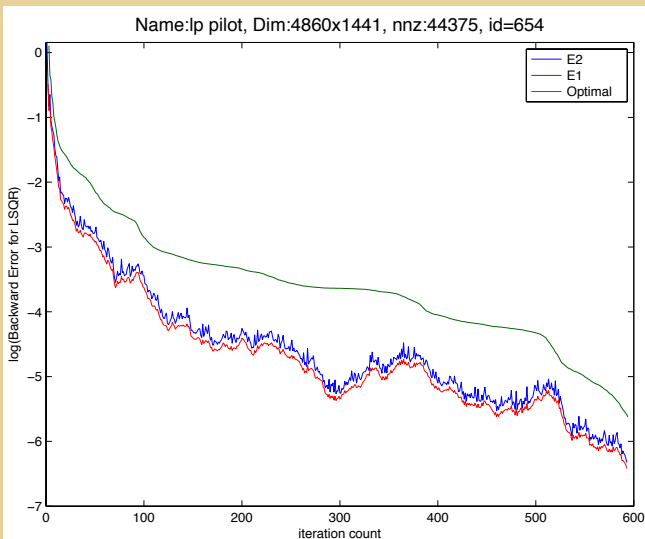
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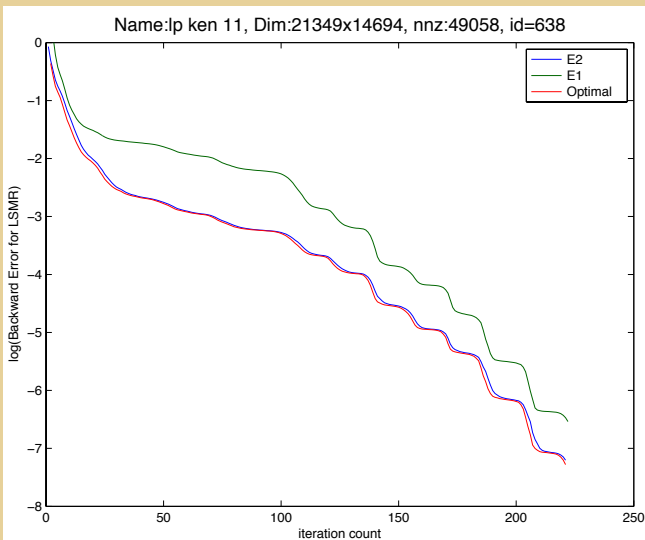
## Backward errors for LSQR – typical



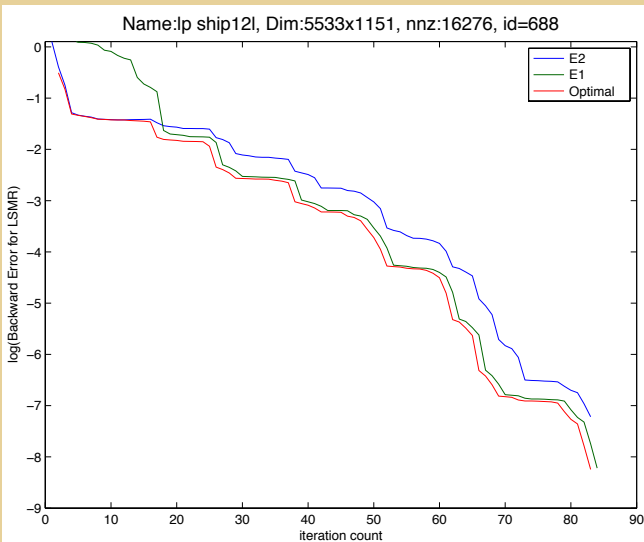
## Backward errors for LSQR – rare



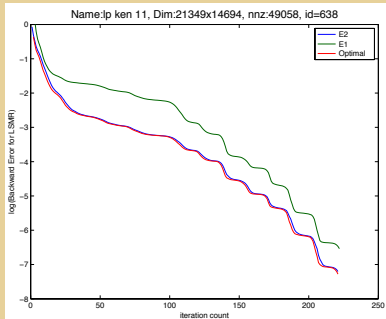
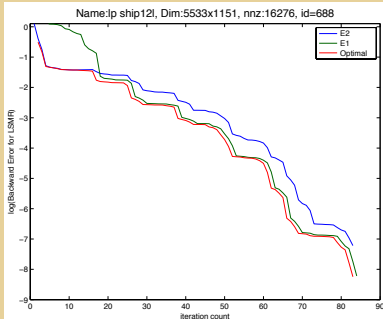
## Backward errors for LSMR – typical



## Backward errors for LSMR – rare



## For LSMR

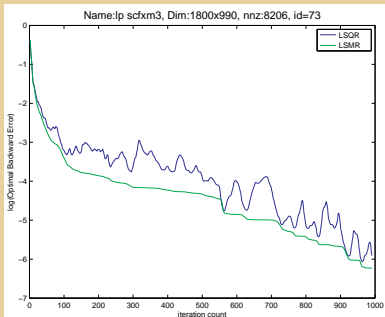
 $\|E_2\| \approx \text{optimal BE almost always}$ 
**Typical:**  $\|E_2\| \approx \tilde{\mu}(x)$ 

**Rare:**  $\|E_1\| \approx \tilde{\mu}(x)$ 


# Optimal backward errors $\tilde{\mu}(x)$

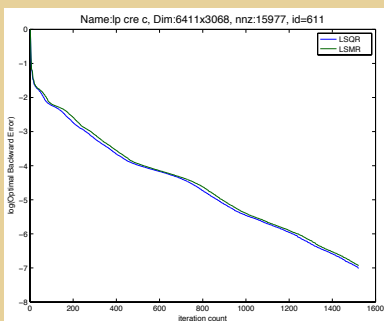
Seem monotonic for LSMR

Usually not for LSQR

## Typical for LSQR and LSMR



## Rare LSQR, typical LSMR

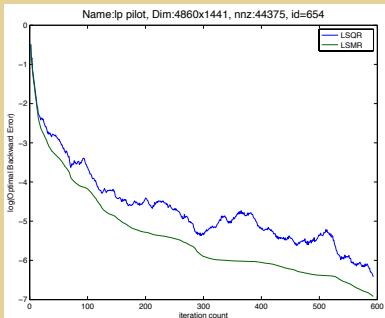




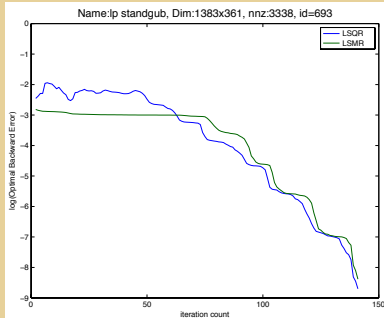
## Optimal backward errors

$$\tilde{\mu}(x^{\text{LSMR}}) \leq \tilde{\mu}(x^{\text{LSQR}}) \text{ almost always}$$

### Typical



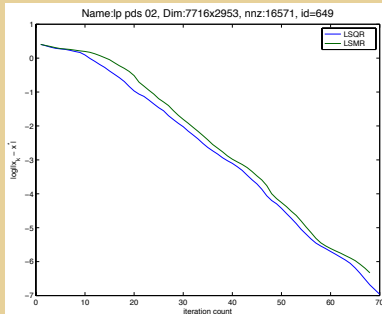
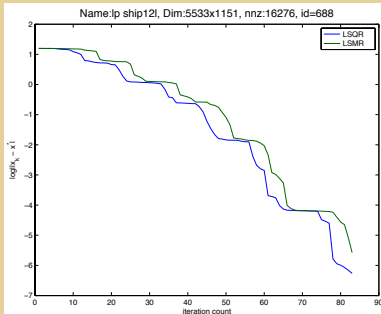
### Rare



# Errors in $x_k$

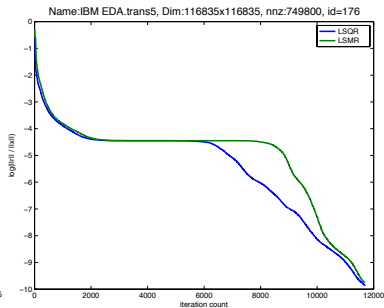
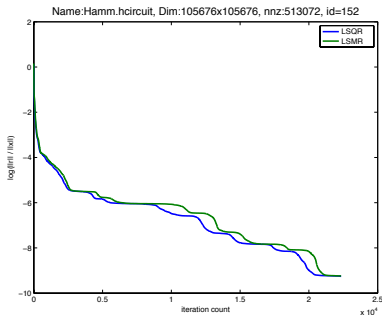
- $\|x^{\text{LSQR}} - x^*\| \leq \|x^{\text{LSMR}} - x^*\|$  seems true

$\|x_k - x^*\|$  for LSMR and LSQR



# Square consistent systems

- $Ax = b$
- Backward error:  $\frac{\|r_k\|}{\|x_k\|}$
- LSQR slightly faster than LSMR in most cases



# Underdetermined systems

Infinitely many solutions

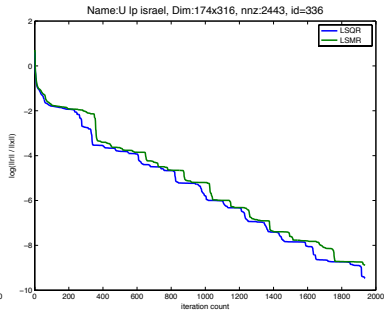
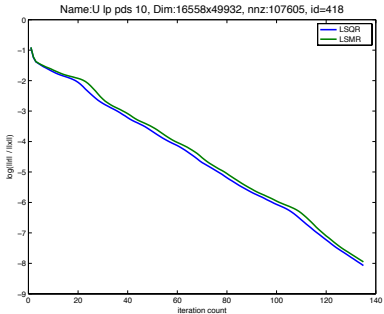
$$Ax = b$$

Unique solution

$$\min \|x\| \text{ st } Ax = b$$

## Theorem

LSQR and LSMR both return the minimum-norm solution



# Summary



# Theoretical properties for $Ax = b, A \succ 0$

## CG and MINRES

$$\begin{array}{ll} \|x^* - x_k\| & \searrow \\ \|x^* - x_k\|_A & \searrow \\ \|x_k\| & \nearrow \end{array}$$



# Theoretical properties for $Ax = b, A \succ 0$

## CG and MINRES

$$\begin{array}{l} \|x^* - x_k\| \quad \searrow \\ \|x^* - x_k\|_A \quad \searrow \\ \|x_k\| \quad \nearrow \end{array}$$

## MINRES

$$\begin{array}{l} \|r_k\| \quad \searrow \\ \|r_k\| / \|x_k\| \quad \checkmark \\ \|r_k\| / (\alpha \|A\| \|x_k\| + \beta \|b\|) \quad \checkmark \end{array}$$

# Theoretical properties for $Ax = b, A \succ 0$

## CG and MINRES

$$\begin{array}{l} \|x^* - x_k\| \quad \searrow \\ \|x^* - x_k\|_A \quad \searrow \\ \|x_k\| \quad \nearrow \end{array}$$

## MINRES

$$\begin{array}{l} \|r_k\| \quad \searrow \\ \|r_k\| / \|x_k\| \quad \checkmark \\ \|r_k\| / (\alpha \|A\| \|x_k\| + \beta \|b\|) \quad \checkmark \end{array}$$

For MINRES, backward errors are monotonic  
 $\Rightarrow$  safe to stop early



# Theoretical properties for $\min \|Ax - b\|$

## LSQR and LSMR

$$\|x^* - x_k\| \quad \searrow$$

$$\|r^* - r_k\| \quad \searrow$$

$$\|r_k\| \quad \searrow$$

$$\|x_k\| \quad \nearrow$$

$x_k \rightarrow$  min-length  $x^*$  if  $\text{rank}(A) < n$



# Theoretical properties for $\min \|Ax - b\|$

## LSQR and LSMR

$$\|x^* - x_k\| \quad \searrow$$

$$\|r^* - r_k\| \quad \searrow$$

$$\|r_k\| \quad \searrow$$

$$\|x_k\| \quad \nearrow$$

$x_k \rightarrow$  min-length  $x^*$  if  $\text{rank}(A) < n$

## LSMR

$$\|A^T r_k\| \quad \searrow$$

$$\|A^T r_k\| / \|r_k\| \quad \searrow$$

almost always

$\approx$  optimal BE almost always

# Theoretical properties for $\min \|Ax - b\|$

## LSQR and LSMR

$$\|x^* - x_k\| \quad \searrow$$

$$\|r^* - r_k\| \quad \searrow$$

$$\|r_k\| \quad \searrow$$

$$\|x_k\| \quad \nearrow$$

$x_k \rightarrow$  min-length  $x^*$  if  $\text{rank}(A) < n$

## LSMR

$$\|A^T r_k\| \quad \searrow$$

$$\|A^T r_k\| / \|r_k\| \quad \searrow$$

almost always

$\approx$  optimal BE almost always

For LSMR, optimal backward errors **seem** monotonic  
 $\Rightarrow$  safe to stop early

## References:

- [LSMR: An iterative algorithm for sparse least-squares problems](#)  
David Fong and Michael Saunders, SISC 2011
- [CG versus MINRES: An empirical comparison](#)  
David Fong and Michael Saunders, SQU Journal for Science 2012

## Kindest thanks:

[Georg Bock and colleagues](#)  
[Phan Thanh An and colleagues](#)

