#### Generalized MINRES and LSQR

#### Orthogonal tridiagonalization of general matrices

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CME 510 Linear Algebra and Optimization Seminar Stanford University, October 3, 2007



- History
- 2 Tridiagonalization of symmetric A

- History
- Tridiagonalization of symmetric A
- $\odot$  Bidiagonalization of rectangular A

- Mistory
- Tridiagonalization of symmetric A
- Bidiagonalization of rectangular A
- 4 Tridiagonalization of unsymmetric A

- History
- Tridiagonalization of symmetric A
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- Mistory
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- **1** Unsymmetric Ax = b

- History
- Tridiagonalization of symmetric A
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- 4 Tridiagonalization of unsymmetric A
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- **7** Unsymmetric Ax = b
- 8 Elizabeth Yip's aim

- History
- Tridiagonalization of symmetric A
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- 4 Tridiagonalization of unsymmetric A
- Original motivation
- 6 Symmetric Ax = b
- **7** Unsymmetric Ax = b
- 8 Elizabeth Yip's aim
- Numerical results

- Mistory
- Tridiagonalization of symmetric A
- 3 Bidiagonalization of rectangular A
- 4 Tridiagonalization of unsymmetric A
- **5** Original motivation
- 6 Symmetric Ax = b
- **7** Unsymmetric Ax = b
- 8 Elizabeth Yip's aim
- Mumerical results
- Conclusions

# History of iterative solvers

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- 1982 Paige-Saunders LSQR Golub-Kahan bidiagonalization for general Ax = b, min ||Ax b||

# History (contd)

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Unsymmetric tridiagonalization, focused on rectangular A

• 2007 Golub, Stoll, and Wathen (draft)

"Approximation of outputs"

Unsymmetric tridiagonalization, focused on Ax = b,  $A^Ty = c$  and estimation of  $c^Tx$  and  $b^Ty$ 

# Tridiagonalization of symmetric *A* using orthogonal matrices

#### Symmetric A

• Tridiagonalization for dense EVD (eigenvalues)

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• Symmetric Lanczos process on A, b

$$\beta_{1}v_{1} = b$$

$$p_{1} = Av_{1} \qquad \alpha_{1} = v_{1}^{T}p_{1}$$

$$\beta_{2}v_{2} = p_{1} - \alpha_{1}v_{1}$$

$$T_{k} = \begin{pmatrix} \alpha_{1} & \beta_{2} & & \\ \beta_{2} & \alpha_{2} & \beta_{3} & & \\ & * & * & * \\ & & \beta_{k} & \alpha_{k} \end{pmatrix}$$

$$p_{2} = Av_{2} \qquad \alpha_{2} = v_{2}^{T}p_{2}$$

$$\beta_{3}v_{3} = p_{2} - \alpha_{2}v_{2} - \beta_{1}v_{1}$$

$$V_{k} = \begin{pmatrix} v_{1} & v_{2} & & \\ & v_{2} & \cdots & v_{k} \end{pmatrix}$$

$$AV_{k} = V_{k}T_{k} + \beta_{k+1}V_{k+1}e_{k}^{T}$$

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# Bidiagonalization of rectangular A

#### Rectangular A

• Bidiagonalization for dense SVD (Golub and Kahan 1965)

$$U^{\mathsf{T}}AV = B \quad \Rightarrow \quad AV = UB, \quad A^{\mathsf{T}}U = VB^{\mathsf{T}}$$

#### Rectangular A

• Bidiagonalization for dense SVD (Golub and Kahan 1965)

• Golub-Kahan process on A, b

$$\beta_1 u_1 = b, \quad \alpha_1 v_1 = A^T u_1$$

$$\beta_2 u_2 = A v_1 - \alpha_1 v_1$$

$$\alpha_2 v_2 = A^T u_2 - \beta_2 v_1$$

$$B_{k} = \begin{pmatrix} \alpha_{1} & & & & \\ \beta_{2} & \alpha_{2} & & & \\ & * & * & & \\ & & \beta_{k} & \alpha_{k} \\ & & & \beta_{k+1} \end{pmatrix}$$

$$U_{k} = \begin{pmatrix} u_{1} & u_{2} & \dots & u_{k} \end{pmatrix}$$

$$V_{k} = \begin{pmatrix} v_{1} & v_{2} & \dots & v_{k} \end{pmatrix}$$

$$AV_k = U_{k+1}B_k, \quad A^TU_k = V_kL_k^T$$

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# Upper or lower bidiagonal?

• Dense A

$$AV = UB = U \begin{pmatrix} * & * & \\ & * & * \\ & & * & * \\ & & & * \end{pmatrix}$$

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• Sparse A with  $b = \beta_1 u_1$ 

$$AV_{k} = U_{k+1}B_{k} \quad \Rightarrow \quad (b \quad AV_{k}) \quad = \quad U_{k+1} \begin{pmatrix} \beta_{1}e_{1} & B_{k} \end{pmatrix}$$

$$\Rightarrow \quad (b \quad A) \begin{pmatrix} 1 & & & \\ & V_{k} \end{pmatrix} \quad = \quad U_{k+1} \begin{pmatrix} * & * & & \\ & * & * & \\ & & * & * \end{pmatrix}$$

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# Tridiagonalization of unsymmetric or rectangular *A* (the "new method")

#### Rectangular A

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Tridiagonalization for dense EVD (eigenvalues)

Bi-tridiagonalization process on A, b, c

$$\beta_{1}u_{1} = b \qquad \gamma_{1}v_{1} = c$$

$$p_{1} = Av_{1} \qquad \alpha_{1} = u_{1}^{T}p_{1}$$

$$\beta_{2}u_{2} = p_{1} - \alpha_{1}u_{1} - \gamma_{1}u_{0}$$

$$T_{k} = \begin{pmatrix} \alpha_{1} & \gamma_{2} & & \\ \beta_{2} & \alpha_{2} & \gamma_{3} & & \\ & * & * & * \\ & & \beta_{k} & \alpha_{k} \end{pmatrix}$$

$$q_{1} = A^{T}u_{2}$$

$$V_{k} = \begin{pmatrix} u_{1} & u_{2} & \dots & u_{k} \\ v_{1} & v_{2} & \dots & v_{k} \end{pmatrix}$$

 $AV_k = U_k T_k + \beta_{k+1} u_{k+1} e_k^T$ 

 $A^T U_{\iota} = V_k T_{\iota}^T + \gamma_{k+1} v_{k+1} e_k^T$ Oct 3, 2007

History Tridiag 1 Bidiag Tridiag 2 Original aim Symmetric Ax = b Unsymmetric Ax = b Yip's aim Results Conclusions

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- Bi-tridiagonalization of unsymmetric A is no more than twice the work and storage per iteration
- If A is symmetric, we get Lanczos
- If A is nearly symmetric, total itns should be not much more

# **Solving symmetric** Ax = b **via Lanczos**

#### Lanzcos process:

$$AV_k = V_{k+1}H_k, \qquad H_k = \begin{pmatrix} \alpha_1 & \beta_2 \\ \beta_2 & \alpha_2 & \beta_3 \\ & * & * & * \\ & & \beta_k & \alpha_k \\ & & & \beta_{k+1} \end{pmatrix}$$

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  $r_k = b - Ax_k$ 

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Suppose  $x_k = V_k w_k$  for some  $w_k$ 

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Three subproblems make  $H_k w_k \approx \beta_1 e_1 \Rightarrow \text{CG, SYMMLQ, MINRES}$  (e.g.  $T_k w_k = \beta_1 e_1$  for CG)

## Symmetric → Unsymmetric

Lanczos on 
$$\begin{pmatrix} I & A \\ A^T & \end{pmatrix} \begin{pmatrix} r \\ x \end{pmatrix} = \begin{pmatrix} b \\ 0 \end{pmatrix}$$
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 (square  $A$ ) is not equivalent to bi-tridiagonalization (but seems worth trying!)

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# **Solving unsymmetric** Ax = b **via bi-tridiagonalization**

#### Bi-tridiag process:

$$AV_k = U_k T_k + \beta_{k+1} u_{k+1} e_k^T \equiv U_{k+1} H_k^{\beta}$$
  

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Similarly, let  $y_k = U_k \bar{w}_k$  to solve  $A^T y = c$ Three subproblems make  $H_k^{\gamma} y_k \approx \gamma_1 e_1$ 

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Not much extra effort to get both  $x_k$  and  $y_k$ 

History Tridiag 1 Bidiag Tridiag 2 Original aim Symmetric Ax = b Unsymmetric Ax = b Yip's aim Results Conclusions

# Elizabeth Yip's motivation (1982)

(Boeing Computer Services Co.)

CG method for unsymmetric matrices applied to PDE problems

We present a CG-type method to solve Ax = b, where A is an arbitrary nonsingular unsymmetric matrix. The algorithm is equivalent to an orthogonal tridiagonalization of A.

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We apply a preconditioned version (Fast Poisson) to the difference equation of unsteady transonic flow with small disturbances. (Compared with ORTHOMIN(5))

# Numerical results with unsymmetric tridiagonalization

## Numerical results (SSY 1988)

$$A = \begin{pmatrix} B & -I \\ -I & B & -I \\ & \ddots & \ddots & \ddots \\ & & -I & B & -I \\ & & & -I & B \end{pmatrix}$$

$$400 \times 400$$

$$B = \text{tridiag} (-1 - \delta \quad 4 \quad -1 + \delta)$$

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Megaflops to reach  $||r|| \le 10^{-6} ||b||$ :

δ	0.0	0.01	0.1	1.0	10.0	100.0
ORTHOMIN(5)	0.31	0.57	0.75	0.83	2.55	2.11
LSQR	0.28	1.38	1.48	0.80	0.57	0.27
USYMQR	0.30	1.88	1.98	1.41	0.99	0.64

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Bottom line:

ORTHOMIN sometimes good, can fail. LSQR always better than USYMQR

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History Tridiag 1 Bidiag Tridiag 2 Original aim Symmetric Ax = b Unsymmetric Ax = b Yip's aim Results Conclusions

# Numerical results (Reichel and Ye 2006)

Focused on rectangular A and least-squares
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Example 1: We know  $x \approx \text{constant}$ . Choose  $c = \begin{pmatrix} 1 & 1 & \dots & 1 \end{pmatrix}^T$ 



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Example 2 (Star cluster)

•  $256 \times 256$  pixels (n = 65536), 470 stars

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- $256 \times 256$  pixels (n = 65536), 470 stars
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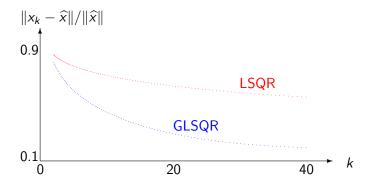
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## **Conclusions**

Slide 27/29

### Subspaces

• Unsymmetric Lanczos generates two Krylov subspaces:

$$U_k \in \operatorname{span}\{b \ Ab \ A^2b \dots A^{k-1}b\}$$
  
 $V_k \in \operatorname{span}\{c \ A^Tc \ (A^T)^2c \dots (A^T)^{k-1}c\}$ 

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Bi-tridiagonalization generates

$$U_{2k} \in \operatorname{span}\{b \ AA^Tb \ \dots \ (AA^T)^{k-1}b \ Ac \ (AA^T)Ac \ \dots\}$$
  
 $V_{2k} \in \operatorname{span}\{c \ A^TAc \ \dots \ (A^TA)^{k-1}c \ A^Tb \ (A^TA)A^Tb \ \dots\}$ 

SOL, Stanford University Oct 3, 2007 Slide 28/29

# Functionals $c^Tx$ , $b^Ty$

• Lu and Darmofal (SISC 2003) use unsymmetric Lanczos with QMR to solve Ax = b and  $A^Ty = c$  simultaneously and to estimate  $c^Tx$  and  $b^Ty$  at a superconvergent rate:

$$|c^T x_k - c^T x| \approx |b^T y_k - b^T y| \approx \frac{\|b - A x_k\| \|c - A^T y_k\|}{\sigma_{\min}(A)}$$

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# Thanks for your patience!!

SOL, Stanford University Oct 3, 2007 Slide 29/29