# Basic Feasible Solutions: A Quick Introduction

CS 261

WILL FOLLOW A CELEBRATED INTELLECTUAL TEACHING TRADITION



"The pland, for high expectations. They plant peet to the lipsol-quart three acress: BARROW ORAMA

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#### TEN STEPS TOWARDS UNDERSTANDING VERTEX OPTIMALITY AND BASIC FEASIBLE SOLUTIONS

THESE SLIDES: MOSTLY INTUITION; PROOFS OMITTED

# Step 0: Notation

- Assume an LP in the following form
  - Maximize  $\mathbf{c}^{\mathsf{T}}\mathbf{x}$
  - Subject to:

 $\mathbf{A}\mathbf{x} \le \mathbf{b}$  $\mathbf{x} \ge \mathbf{0}$ 

- N Variables, M constraints
- U = Set of all feasible solutions

### Step 1: Convex Sets and Convex Combinations

 Convex set: If two points belong to the set, then any point on the line segment joining them also belongs to the set

- Convex combination: weighted average of two or more points, such that the sum of weights is 1 and all weights are non-negative
  - Simple example: average

#### **Convex Combination**

- Given points x<sup>1</sup>, x<sup>2</sup>, ..., x<sup>K</sup>, in N dimensions
- For any scalars a<sub>1</sub>, a<sub>2</sub>, ..., a<sub>K</sub> such that
  Each a<sub>i</sub> is non-negative

$$-a_1 + a_2 + \dots + a_K = 1$$

 $a_1x^1 + a_2x^2 + ... + a_Kx^K$  is a convex combination of  $x^1, x^2, ..., x^K$ 

# Example: Convex combination of two points



#### Exercise

 If w is a convex combination of x and y, and z is a convex combination of x and w, then must z be a convex combination of x and y?

# Example: Convex combination of three points



#### Exercise

Is the set of all convex combinations of N points always convex?



Can any convex set be represented as a convex combination of N points, for finite N?

#### Step 2: Half spaces are convex

•  $\mathbf{a}^{\mathsf{T}}\mathbf{x} \leq \mathbf{b}, \, \mathbf{a}^{\mathsf{T}}\mathbf{y} \leq \mathbf{b}$ 

- Choose **z** = (**x**+**y**)/2
- Then, we must have  $\mathbf{a}^{\mathsf{T}}\mathbf{z} \leq \mathbf{b}$

# Step 3: Intersection of convex sets is convex

- Two convex sets S, T
- Suppose **x**, **y** both belong to  $S \cap T$
- z = (x+y)/2
  - (1) z belongs to S since S is convex
  - (2) **z** belongs to T since T is convex
- Hence, **z** belongs to  $S \cap T$
- Implication: S∩T is convex

# Step 4: Feasible solutions to LP

 The set U of feasible solutions to a Linear Program is convex

• Proof: U is the intersection of Half Spaces

#### Exercise

• Must every convex set be bounded?

### Step 5: Basic Feasible Solution

• **x** is a basic feasible solution to a LP, if

(1) x is a feasible solution

(2) There do not exist two other feasible solutions y, z such that x = (y+z)/2

ALSO known as vertex solution, extreme point solution, corner-point solution



# Step 6: BFS and Bounded Polytopes

 If U is bounded, then any point in U is the convex combination of its basic feasible solutions

Use  $\mathbf{b}_1$ ,  $\mathbf{b}_2$ , ...,  $\mathbf{b}_K$  to denote the K basic feasible solutions



# Step 7: BFS and Optimality for Bounded Polytopes

 If U is bounded and non-empty, then there exists an optimal solution that is also basic feasible

### Step 8: BFS and Optimality for General LPs

 If an LP has a basic feasible solution and an optimum solution, then there exists an optimal solution that is also basic feasible

 Will call these solutions "Vertex Optimal" or "Optimum Basic Feasible"

#### **Step 9: Two Implications**

(1) SOLUTION OF LPs: The Simplex method which walks from bfs to bfs is correct for bounded polytopes

(2) STRUCTURAL CONSEQUENCES: The Vertex Optimal Solutions often have very interesting properties. Example: For the matching LP in the next video, every vertex optimal solution is integral



## In Practice

- Most LP Solvers return an optimum basic feasible solution, when one exists.
  - Either, they use Simplex
  - Or, they transform the solution that they do find to a basic feasible solution
- Hence, when we solve a problem using Excel we get an optimum basic feasible solution, when one exists.

# Step 10: A useful property

- At every basic feasible solution, at least N constraints are tight
- More specifically: Exactly N Linearly Independent Constraints are tight
- Basic solution: Not necessarily feasible, but exactly N Linearly Independent Constraints are tight



# Is it possible for a LP to have an optimum solution but no basic feasible solution?



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#### Exercise

You are given a LP with three decision variables where some of the constraints are

 $0 \le x \le 1$  $0 \le y \le 1$  $0 \le z \le 1$ 

Is it possible to add other constraints and an objective function such that the LP has an optimum solution but not an optimum basic feasible solution?

# LP in standard form

Maximize  $\mathbf{c}^{\mathsf{T}}\mathbf{x}$ 

Subject to:

$$Ax = b$$

**x** ≥ 0

- N Variables, M constraints, assume that the M constraints are Linearly Independent
- At any BFS, at most M variables are strictly positive

#### Recap

- A feasible solution is basic feasible if it is not the average of two other feasible solutions
- If the feasibility region U for a LP is bounded and nonempty, then there exists an optimal solution that is also basic feasible
- If an LP has a basic feasible solution and an optimum solution, then there exists an optimal solution that is also basic feasible
- Commercial LP solvers return an optimum basic feasible solution by default, where one exists
- If there are N decision variables, then a basic feasible solution has N linearly independent constraints tight