

Linear Recurrences

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$$\{X_n\} : X_{n+k} = a_{k-1} X_{n+k-1} + \dots + a_0 X_n$$

① Linearity (of sequences)

$$X_{n+k} = a_{k-1} X_{n+k-1} + \dots + a_0 X_n \quad \forall$$

$$y_{n+k} = a_{k-1} y_{n+k-1} + \dots + a_0 y_n$$

$$(X_{n+k} + y_{n+k}) = a_{k-1} (X_{n+k-1} + y_{n+k-1}) + \dots + a_0 (X_n + y_n)$$

$\{X_n\}, \{y_n\}$ satisfy linearity, so will the sum $\{X_n + y_n\}$

② $k=1$. $X_{n+1} = r X_n$. Sol: $X_n = A \cdot r^n$

When is this a solution for linearity?

$$r^k = a_{k-1} r^{k-1} + \dots + a_0$$

(characteristic
(LHS - RHS)
poly - in
"characteristic"

$$X_n = A_1 r_1^n + \dots + A_m r_m^n$$

are solutions!

^ - , , , , ^ -

(3) Elementary way to solve this.

E.g. $X_{n+2} = 5X_{n+1} - 6X_n$. $[x^2 - 5x + 6 = 0]$
 $x = 2$ is a root
 $\times (1 - 3x)?$
 $Y_{n+1} = 2Y_n$
 $Y_n = A \cdot 2^n$

$$X_{n+2} - 3X_{n+1} = A \cdot 2^n$$

$$X'_{n+2} - 3X'_{n+1} = A \cdot 2^n$$

$$Z_{n+2} - 3Z_{n+1} = 0$$

$$X'_n = \left(-\frac{1}{2}\right) A \cdot 2^n$$

$$Z_n = X_n - X'_n = B \cdot 3^n$$

Think about

$$X_{n+2} = 4X_{n+1} - 4X_n, \quad (x-2)^2 = 0$$

$$X_n = (A_n + B_n) \cdot 2^n$$

(4) Finite differences.

char Eq: $(x-1)^k$ $\{X_n\} \xrightarrow{\Delta} \{X_n - X_{n-1}\}$

$\Leftrightarrow \Delta^k \{X_n\} = \{0\}$ iff $\deg k$ poly $\deg k$ poly

$\Leftrightarrow X_n$ is a $\deg k-1$ poly $(in\ n)$ $k \geq 1$

$\{X_n\}$ 1, 4, 9, 16, 25, 36

$$\Delta \{x_n\} \quad \Delta^2 \{x_n\}$$

(5) "Continuation" (\square is "deg k")

x_0, \dots, x_{k-1} known, Then rest of seq. is
Can continue this backwards. x_{-1} is well-defined

Think about: \mathbb{Z} , $a_i \in \mathbb{Z}$

(6) Main Thm.

Fact (General Formula for Linear Recurrences)

Suppose $\{x_i\}$ is a sequence satisfying

$$x_{n+k} = a_{k-1}x_{n+k-1} + a_{k-2}x_{n+k-2} + \dots + a_0x_n$$

for all integers $n \geq 0$. Then x_n has the following general formula:

$$x_n = P_1(n)\alpha_1^n + P_2(n)\alpha_2^n + \dots + P_m(n)\alpha_m^n$$

where the α_i are the roots of the **characteristic polynomial**:

$$x^k - a_{k-1}x^{k-1} - a_{k-2}x^{k-2} - \dots - a_0 = (x - \alpha_1)^{\beta_1}(x - \alpha_2)^{\beta_2} \dots (x - \alpha_m)^{\beta_m}$$

and P_i are polynomials with degree at most $\beta_i - 1$

$$(x-2)^2(x+1)$$

$$x_n = A \cdot 2^n + B \cdot n 2^n + C \cdot (-1)^n$$

(7) Some Tricks

$$\sum a_i x_i^n$$

Rec \longleftrightarrow Sum of Exp.

$$(x-x)(x-y)=0$$

$$ax^n + by^n$$

$$x_{n+1} = (x+y)x_n - (x-y)x_{n-1}$$

- Linearity is so so so imp!

$$\{X_n - \lambda X_{n-1}\}$$

- Rec \longleftrightarrow Combi

Ex $a_1 = 1, \sum_{d|n} a_d = 2^n$

Show $n | a_n$.



Hints

Q1 $\alpha > 1, \lfloor \alpha^n \rfloor, \frac{1}{\alpha^n} < 1$

$$\underbrace{\alpha^n + \frac{1}{\alpha^n}}_{\text{force to } \equiv 1 \pmod{p}} = \lfloor \alpha^n \rfloor + 1 \quad \left| \begin{array}{l} n \neq 0 \\ 2 \not\equiv 1 \pmod{p} \end{array} \right.$$

Q2 $f(0) = 0, f \text{ odd. } \boxed{\checkmark}$

$a \times b \times c = 0$ is strong.

$$c = 0$$

f non-decr.

$$x > 0$$

$$\dots f\left(\frac{x}{i}\right), f(x), f(2x), f(4x), \dots \geq 0$$

C_i not all zero. $\sum C_i V_i = 0$ | "dependent"

Linear algebra machinery.

$$\det \begin{pmatrix} 1 & 1 & \dots & 1 \\ v_1 & v_2 & \dots & v_n \\ 1 & 1 & \dots & 1 \end{pmatrix} \begin{cases} = 0 & \text{if dep.} \\ \neq 0 & \text{if indep.} \end{cases}$$

"poly of γ_i 's with total dg. $\frac{n(n-1)}{2}$ "

$$\begin{pmatrix} \gamma_1 & \gamma_2 & \dots & \gamma_k \\ \vdots & \vdots & \dots & \vdots \\ \gamma_1^{k-1} & \gamma_2^{k-1} & \dots & \gamma_k^{k-1} \end{pmatrix} \prod_{i < j} (\gamma_i - \gamma_j)$$

homogeneous.

"Vandermonde matrix"

(9) Analysis (?)

$$\sum C_i \vec{v}_i = 0 \rightarrow \sum C_i \gamma_i^n = 0 \text{ for } n=0, \dots$$

also satisfies the
deg $< \text{recurrence}$.

\Rightarrow holds for all n

Impossible if $\gamma_i \in \mathbb{R}_{>0}$

$$\sum C_i \gamma_i^n = - \sum C_i \gamma_i^n$$

γ_i : max. size
 $|\gamma_i| = M$

γ_j : others
 $|\gamma_j| < M$

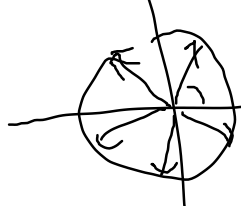
$$\sum C_i \left(\frac{\gamma_i}{M} \right)^n = - \sum C_j \left(\frac{\gamma_j}{M} \right)^n$$

" $\rightarrow 0$
 $1 \cdot 1 = 1$

Pick N larger s.t.
 and larger

$\left(\frac{\gamma_i}{M} \right)^N \approx 1$

$\nexists |\gamma_i| = M$ sim



$$\sum C_i \left(\frac{\gamma_i}{M} \right)^N \approx \sum C_i = 0$$

\downarrow
 0

$$\sum (C_i \gamma_i) \gamma_i^N \rightarrow \sum_{|\gamma_i| \max} C_i \gamma_i = 0$$

$$\therefore \sum_{|\gamma_i| \max} C_i \gamma_i^N = 0 \quad \forall N \Rightarrow C_i = 0 \quad (?)$$

(1) Limits of recursive relations

rec. rel. \longleftrightarrow characteristic poly

$$\{X_n\} : P, Q \Rightarrow P+Q$$

$$\underline{P \cdot R}, R_a$$

Euclidean deg!

"minimal char. poly"

$$A_n = \gamma_1^n + \gamma_2^n + \dots + \gamma_k^n$$

$$(\gamma - \gamma_1)(\gamma - \gamma_2) \dots (\gamma - \gamma_k)$$

$$\vec{A}_0 =$$

$$\vec{A}_1 =$$

$$\sum c_i \vec{v}_i \Rightarrow \sum c_i \vec{A}_i \text{ can cover everything}$$

$$\vec{A}_i = (A_{i,1}, A_{i,2}, \dots, A_{i,n})^T$$

$$\sum c_i \vec{A}_i = 0 \Rightarrow \sum c_i A_{i+n} = 0$$

for $n=0, \dots$

$$\Rightarrow \sum_{i=1}^k c_i A_{i+n} = 0 \quad \forall$$

rec. rel. of deg $\leq k$

root of unith filter

Ref: Evan Chen "Napkin"

Also,

Recurrence relations \approx Matrix multiplication

$$\left(\begin{array}{c|cccc} 0 & 1 & & & \\ & & 1 & & \\ & & & \ddots & \\ \hline a_0 & a_1 & \dots & a_{k-1} & \end{array} \right) \begin{pmatrix} x_n \\ \vdots \\ x_{n+k-1} \end{pmatrix} = \begin{pmatrix} x_{n+1} \\ \vdots \\ x_{n+k} \end{pmatrix}$$

γ_i - eigenvalues

More hints

Q5 "Indicator function".

$$F_n = \begin{pmatrix} 0 \text{ if } n \text{ is} \\ \text{in AP \#1} \end{pmatrix} \begin{pmatrix} 0 \text{ if } n \text{ is} \\ \text{in AP \#2} \end{pmatrix}$$

Q2 Find another a s.t. $f(ax) = af(x)$.
(or many other)

Q4 $a_n = \gamma_1^n + \dots + \gamma_k^n$

B Get to ① first!

Q3 $\{f(n)\}$ is a recurrence (in n)
What else?