

Context-free Grammar and Backtrack Parsing

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Problems with Regular Language

Is English a regular language?

Bad question! We do not even know what English is!

Two eggs and bacon make(s) a big breakfast

Can you slide me the salt?

He didn't ought to do that

But—No!

I put the wine you brought in the fridge

I put the wine you brought for Sandy in the fridge

Should we bring the wine you put in the fridge out
now?

You said you thought nobody had the right to claim
that they were above the law

Problems with Regular Language

You said you thought nobody had the right to claim
that they were above the law

Problems with Regular Language

[You said you thought [nobody had the right [to claim
that [they were above the law]]]]

Problems with Regular Language

Is English morphology a regular language?

Bad question! We do not even know what English morphology is!

They sell collectables of all sorts

This concerns unredecontaminability

This really is an untiable knot.

But—Probably!

(Not sure about Swahili, though)

Context-free Grammar

Nonterminal symbols ~ grammatical categories

Terminal Symbols ~ words

Productions ~ (unordered) (rewriting) rules

Distinguished Symbol

Not all that important

- Terminals and nonterminals are disjoint
- Distinguished symbol

Context-free Grammar

Nonterminal symbols ~ grammatical categories

S, NP, VP, Det, N, V

Terminal Symbols ~ words

the, dog, cat, chased

Productions ~ (unordered) (rewriting) rules

S → NP VP	Det → the
NP → Det N	N → dog
VP → V NP	N → cat
	V → chased

Distinguished Symbol

S

Context-free Grammar

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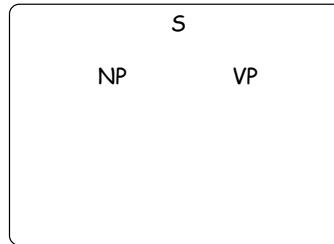
the, dog, cat, chased

Productions ~ (unordered) (rewriting) rules

S → NP VP
 NP → Det N
 VP → V NP
 Det → the
 N → dog
 N → cat
 V → chased

Distinguished Symbol

S



Context-free Grammar

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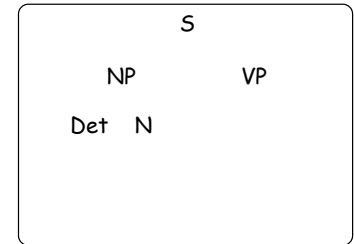
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Distinguished Symbol

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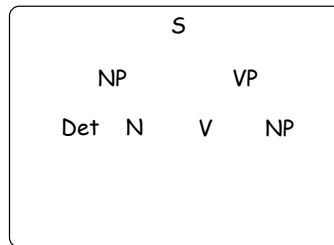
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Productions ~ (unordered) (rewriting) rules

S → NP VP
 NP → Det N
VP → V NP
 Det → the
 N → dog
 N → cat
 V → chased

Distinguished Symbol

S



Context-free Grammar

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Terminal Symbols ~ words

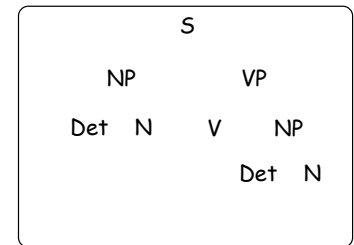
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S → NP VP
NP → Det N
 VP → V NP
 Det → the
 N → dog
 N → cat
 V → chased

Distinguished Symbol

S



Context-free Grammar

Nonterminal symbols ~ grammatical categories

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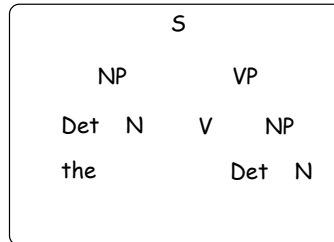
the, dog, cat, chased

Productions ~ (unordered) (rewriting) rules

$S \rightarrow NP VP$
 $NP \rightarrow Det N$
 $VP \rightarrow V NP$
 $Det \rightarrow the$
 $N \rightarrow dog$
 $N \rightarrow cat$
 $V \rightarrow chased$

Distinguished Symbol

S



Context-free Grammar

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Terminal Symbols ~ words

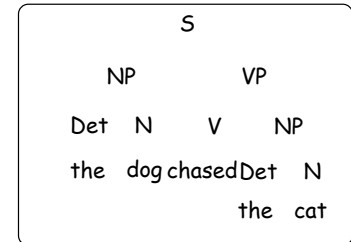
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Context-free Grammar

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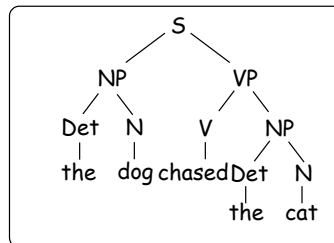
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$S \rightarrow NP VP$
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Distinguished Symbol

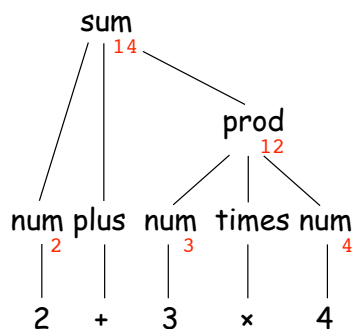
S



Context-free Grammar

- Defines a language (set of strings)
- and
- Defines a corresponding set of structures

Structure and Semantics



Aithmetic Expressions

sum \Rightarrow prod add_op sum

sum \Rightarrow prod

prod \Rightarrow base mul_op prod

prod \Rightarrow base

base \Rightarrow a | b | c ...

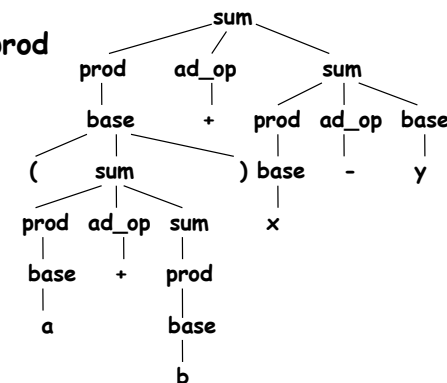
base \Rightarrow (sum)

ad_op \Rightarrow +

ad_op \Rightarrow -

mul_op \Rightarrow *

mul_op \Rightarrow /



Aithmetic Expressions

sum \Rightarrow prod add_op sum

sum \Rightarrow prod

prod \Rightarrow base mul_op prod

prod \Rightarrow base

base \Rightarrow a | b | c ...

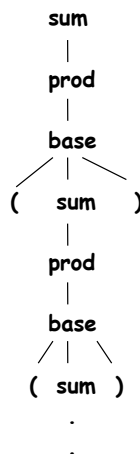
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ad_op \Rightarrow -

mul_op \Rightarrow *

mul_op \Rightarrow /

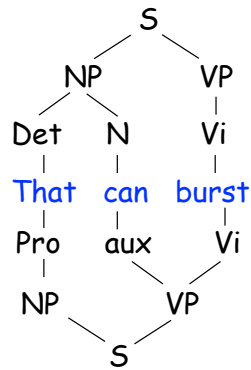


Embedding

Arbitrarily embedded parentheses are not possible in a regular language because the number of open parentheses that still requiring matching close parentheses can be greater than the number of states in the automaton, whatever that number is. Hence, by the pumping lemma, an automaton with only a finite number of states cannot characterize the language.

Ambiguity

More than one structure can correspond to a single string



Lots of Problems!

That it rained slept soundly.

Mary told the boy Mary
John told that it rained
me

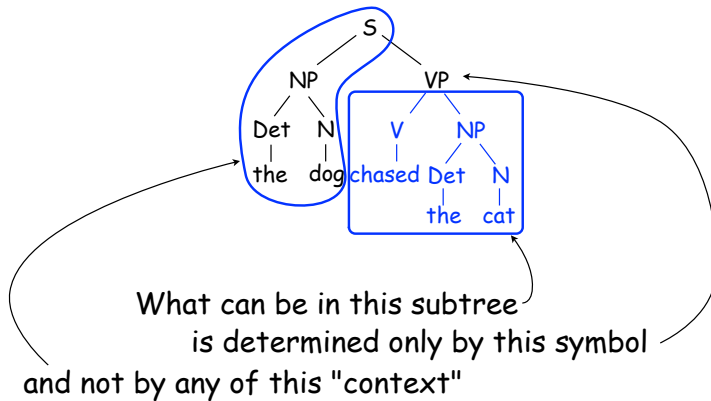
That it rained were certainly true.

His arguments convinces me.
Me told Mary it.

For the moment, just forget the problems

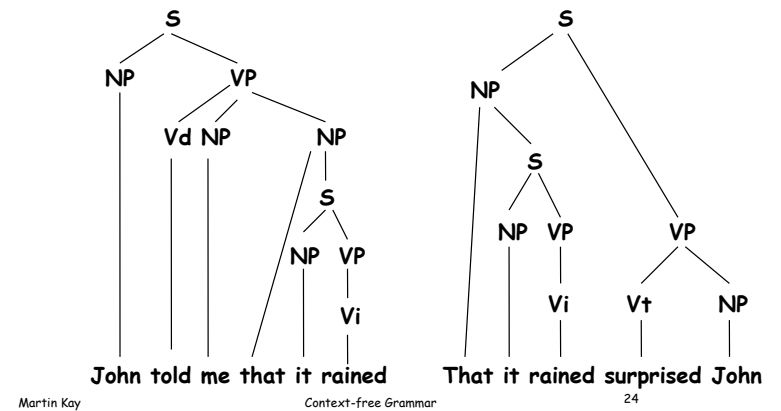


Why "Context-free"?



Convenience

You can define something once, and then use it in many places in the grammar

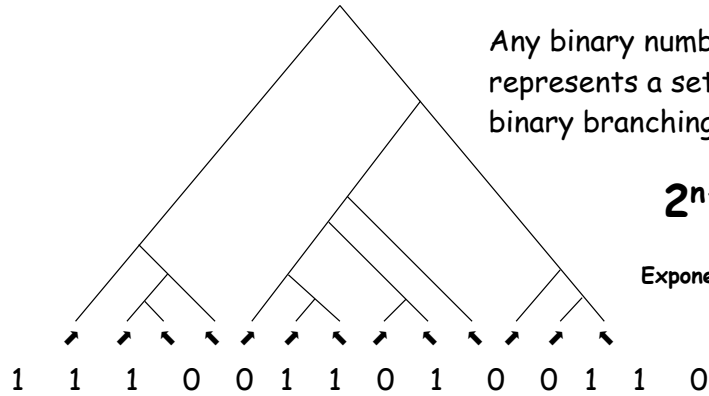


How many (binary) trees with n terminals?

Any binary number represents a set of binary branching trees

$$2^{n-2}$$

Exponential

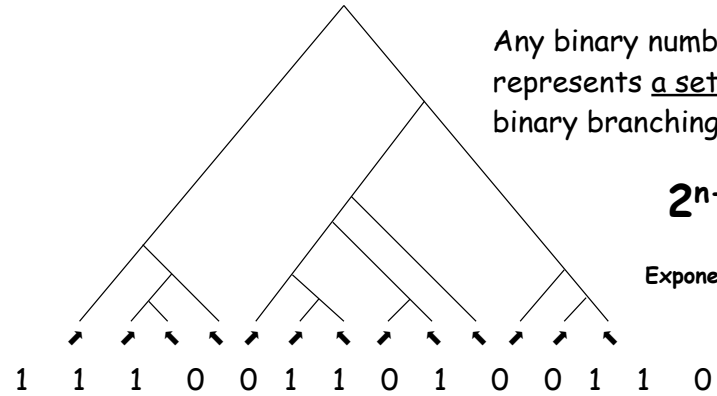


How many (binary) trees with n terminals?

Any binary number represents a set of binary branching trees

$$2^{n-2}$$

Exponential



Catalan Numbers

1 terminal

Terminals Trees

1 1

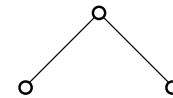
○

Catalan Numbers

2 terminals

Terminals Trees

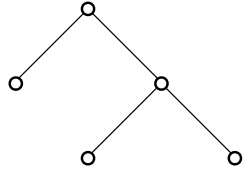
1 1
2 1



Catalan Numbers

3 terminals

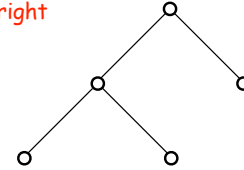
Terminals	Trees
left → 1	1
2	1 ← right



Catalan Numbers

3 terminals

Terminals	Trees
1	1 ← right
left → 2	1

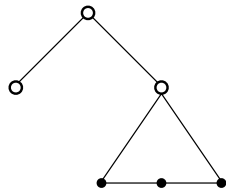


Catalan Numbers

4 terminals

Terminals	Trees
left → 1	1
2	1
3	2 ← right

$$1 \times 2 = 2$$



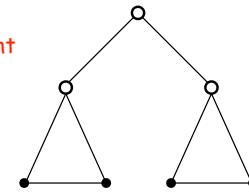
Catalan Numbers

4 terminals

Terminals	Trees
1	1
left → 2	1 ← right
3	2

$$1 \times 2 = 2$$

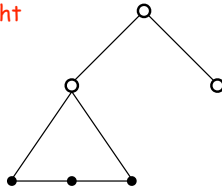
$$1 \times 1 = 1$$



Catalan Numbers

4 terminals

Terminals	Trees
1	1 ← right
2	1
left → 3	2



$$1 \times 2 = 2$$

$$1 \times 1 = 1$$

$$2 \times 1 = 2$$

Catalan Numbers

4 terminals

Terminals	Trees
1	1 ← right
2	1
left → 3	2
4	5

$$1 \times 2 = 2$$

$$1 \times 1 = 1$$

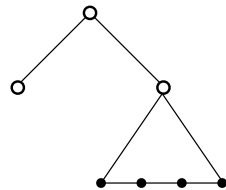
$$\underline{2 \times 1 = 2}$$

$$5$$

Catalan Numbers

5 terminals

Terminals	Trees
left → 1	1
2	1
3	2
4	5 ← right

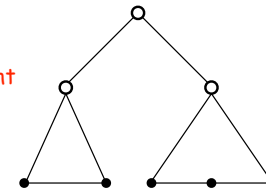


$$1 \times 5 = 5$$

Catalan Numbers

4 terminals

Terminals	Trees
1	1
left → 2	1
3	2 ← right
4	5



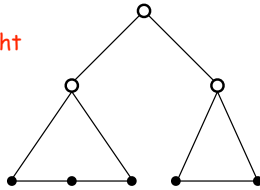
$$1 \times 5 = 5$$

$$1 \times 2 = 2$$

Catalan Numbers

4 terminals

Terminals	Trees
1	1
2	1 ← right
left → 3	2
4	5



$$1 \times 5 = 5$$

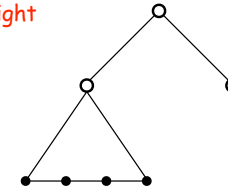
$$1 \times 2 = 2$$

$$2 \times 1 = 2$$

Catalan Numbers

4 terminals

Terminals	Trees
1	1 ← right
2	1
3	2
left → 4	5



$$1 \times 5 = 5$$

$$1 \times 2 = 2$$

$$2 \times 1 = 2$$

$$\underline{5 \times 1 = 5}$$

$$14$$

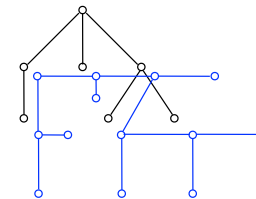
Catalan Numbers

Terminals	Trees
1	1
2	1
3	2 (= 1*1 + 1*1)
4	5 (= 1*2 + 1*1 + 2*1)
5	14 (= 1*5 + 1*2 + 2*1 + 5*1)
6	42 (= 1*14 + 1*5 + 2*2 + 5*1 + 14*1)
...	

$$\text{Cat}(n) = \binom{2n-1}{n}$$

Catalan Numbers

The number of binary trees with n terminals
 The number of arbitrarily branching trees with n nodes.

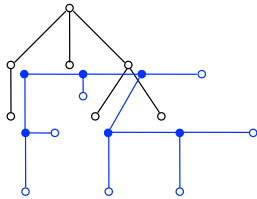


Any tree can be encoded in a binary tree with the same number of terminals as the original tree has nodes.

Catalan Numbers

There are n nonterminals in the **binary tree** for the n nodes in the original tree, except for the root.

There are $n+1$ terminals in a binary tree with n nonterminals. \square



Let there be:

a arcs,
 n nonterminals, and
 t terminals

There are two arcs downward from each nonterminal, so

$$a = 2n$$

There is one arc upward from each terminal and nonterminal except the root, so

$$a = n + t - 1$$

Ergo

$$n = t - 1$$

Analyzing Sentences

- **Recognition:** determining if a given string is a member of the language defined by the grammar
- **Parsing:** Determining what structure(s) a grammar assigns to a string
- **Note:** Parsing encompasses recognition.

Top-down

① the ② dog ③ chased ④ the ⑤ cat ⑥

$S \rightarrow NP VP$ $Det \rightarrow the$

$NP \rightarrow Det N$ $N \rightarrow dog$

$VP \rightarrow V NP$ $N \rightarrow cat$

$V \rightarrow chased$

If there is a sentence between ① and ⑥, then there must be an NP between ① and some \otimes and a VP between \otimes and ⑥.

If there is an NP between ① and \otimes , then there must be a Det between ① and some \heartsuit and a N between \heartsuit and \otimes .

More procedurally

① the ② dog ③ chased ④ the ⑤ cat ⑥

$S \rightarrow NP VP$ $Det \rightarrow the$

$NP \rightarrow Det N$ $N \rightarrow dog$

$VP \rightarrow V NP$ $N \rightarrow cat$

$V \rightarrow chased$

If there is a sentence between ① and ⑥, look for an NP beginning at ① and let \otimes be where it ends.

Look for a VP beginning at \otimes and ending at ⑥.

Definite-clause grammar

```
s(A, C) :- np(A, B), vp(B, C).
vp(A, C) :- v(A, B), np(B, C).
np(A, C) :- det(A, B), n(B, C).
det(A, B) :- the(A, B).
n(A, B) :- dog(A, B).
n(A, B) :- cat(A, B).
v(A, B) :- chased(A, B).
```

```
the(1, 2).
dog(2, 3).
chased(3, 4).
the(4, 5).
cat(5, 6).
```

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Context-free Grammar

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```
| ?- s(X, Y).
X = 1,
Y = 6 ?
| ?- np(X, Y).
X = 1,
Y = 3 ? ;
X = 4,
Y = 6 ? ;
no
| ?-
```

Top-down

These do not have to be numbers. They could be ... anything! Just so they are different from one another.

① the ② dog ③ chased ④ the ⑤ cat ⑥

```
S → NP VP      Det → the
NP → Det N      N → dog
VP → V NP       N → cat
V → chased
```

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Context-free Grammar

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Definite-clause grammar

```
s(A, C) :- np(A, B), vp(B, C).
vp(A, C) :- v(A, B), np(B, C).
np(A, C) :- det(A, B), n(B, C).
det(A, B) :- the(A, B).
n(A, B) :- dog(A, B).
n(A, B) :- cat(A, B).
v(A, B) :- chased(A, B).
```

```
the(start, the1).
dog(the1, dog).
chased(dog, chased)
the(chased, the2).
cat(the2, end).
```

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Context-free Grammar

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```
| ?- s(X, Y).
X = start,
Y = end ? ;
no
| ?- np(X, Y).
X = start,
Y = dog ? ;
X = chased,
Y = end ? ;
no
```

The sentence has to be part of the program!

Top-down

These do not have to be numbers. They could be ... anything! Just so they are different from one another.

① the ② dog ③ chased ④ the ⑤ cat ⑥

```
S → NP VP      Det → the
NP → Det N      N → dog
VP → V NP       N → cat
V → chased
```

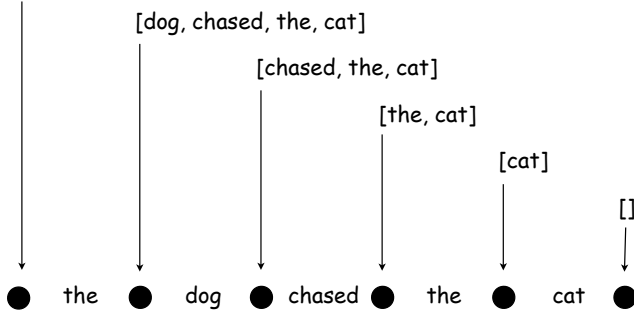
So let's use tails of the string itself!

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Context-free Grammar

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Top-down



If there is a sentence between [the, dog, chased, the, cat] and [], then there must be an NP between [the, dog, chased, the, cat] and some X and a VP between that X and [].

Renamed String Positions

```
s(A, C) :- np(A, B), vp(B, C).
vp(A, C) :- v(A, B), np(B, C).
np(A, C) :- det(A, B), n(B, C).
det([the | X], X).
n([dog | X], X).
n([cat | X], X).
v([chased | X], X).
```

There is a noun between positions A and B if A is a string beginning with the word "dog" and B is the remainder of that string.

Renamed String Positions

```
s(A, C) :- np(A, B), vp(B, C).
vp(A, C) :- v(A, B), np(B, C).
np(A, C) :- det(A, B), n(B, C).
det([the | X], X).
n([dog | X], X).
n([cat | X], X).
v([chased | X], X).
```

| ?- s([the,dog,chased,the,cat], []).
yes

But its only a recognizer!

```
s(A, C, s(NP, VP)) :- np(A, B, NP), vp(B, C, VP).
vp(A, C, vp(V, NP)) :- v(A, B, V), np(B, C, NP).
np(A, C, np(Det, N)) :- det(A, B, Det), n(B, C, N).
det([the | X], X, det(the)).
n([dog | X], X, n(dog)).
n([cat | X], X, n(cat)).
v([chased | X], X, v(chased)).
```

| ?- s([the,dog,chased,the,cat], [], S).
S = s(np(det(the),n(dog)),vp(v(chased),np(det(the),n(cat)))) ?

Caveat Parsor

```
s(A, C, s(NP, VP)) :- np(A, B, NP), vp(B, C, VP).
vp(A, C, vp(V, NP)) :- v(A, B, V), np(B, C, NP).
np(A, C, np(Det, N)) :- det(A, B, Det), n(B, C, N).
det([the | X], X, det(the)).
pp(A, C, pp(P, NP)) :- p(A, B, P), np(B, C, NP).
np(A, C, np(NP, PP)) :- np(A, B, NP), pp(B, C, PP).
vp(A, C, vp(VP, PP)) :- vp(A, B, VP), pp(B, C, PP).
n([dog | X], X, n(dog)).
n([cat | X], X, n(cat)).
| ?- s([the,dog,chased,the,cat], [], S).
n
S = s(np(det(the),n(dog)),vp(v(chased),np(det(the),n(cat)))) ? ;
v
! Resource error: insufficient memory
p([round | X], X, p(round)).
```

Left Recursion

NP --> NP PP

To find an NP, first find an NP !!!

NP --> Det N

Det --> NP apostrophe-s

Left Recursion

A grammar that contains instances of left recursion can be replaced by one with no left recursion which

- describes the same language.
- assigns the same structure to the sentences.

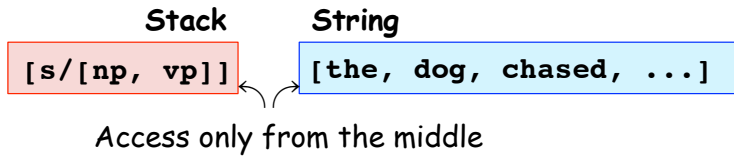
COMING SOON!

For the moment, we will just look for a different parsing algorithm that does not have the problem.

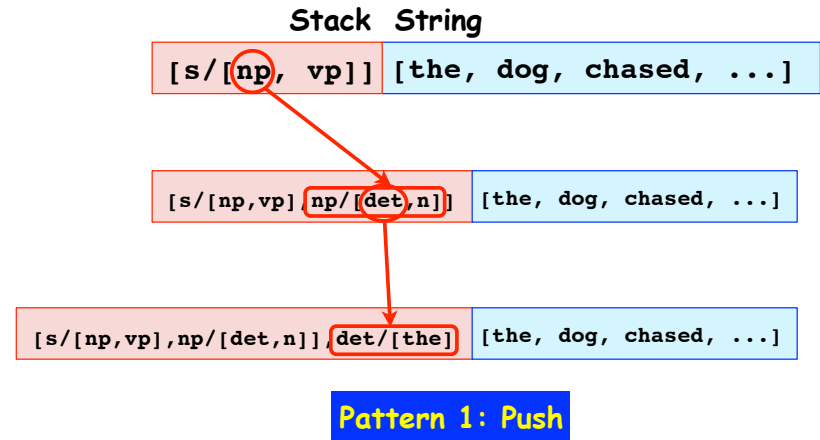


Also, for a while, we will concentrate on recognizers because we know how to turn them into parsers.

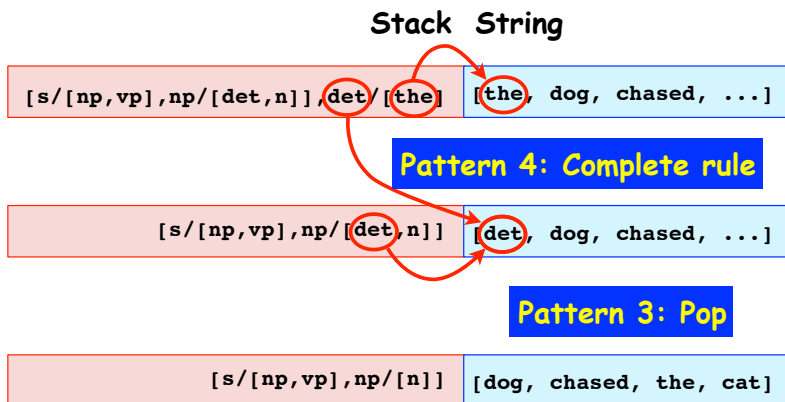
Stack-String Models



Stack-String Models



Stack-String Models



0.	[s/[np, vp]]	[the, dog, chased, the, cat]	0
1.	[s/[np, vp], np/[det, n]]	[the, dog, chased, the, cat]	1
2.	[s/[np, vp], np/[det, n], det/[the]]	[the, dog, chased, the, cat]	1
3.	[s/[np, vp], np/[det, n]]	[det, dog, chased, the, cat]	4
4.	[s/[np, vp], np/[n]]	[dog, chased, the, cat]	3
5.	[s/[np, vp], np/[n], n/[dog]]	[dog, chased, the, cat]	1
6.	[s/[np, vp], np/[n]]	[n, chased, the, cat]	4
7.	[s/[np, vp]]	[np, chased, the, cat]	4
8.	[s/[vp]]	[chased, the, cat]	3
9.	[s/[vp], vp/[v, np]]	[chased, the, cat]	1
10.	[s/[vp], vp/[v, np], v/[chased]]	[chased, the, cat]	1
11.	[s/[vp], vp/[v, np]]	[v, the, cat]	4
12.	[s/[vp], vp/[np]]	[the, cat]	3
13.	[s/[vp], vp/[np], np/[det, n]]	[the, cat]	1
14.	[s/[vp], vp/[np], np/[det, n], det/[the]]	[the, cat]	1
15.	[s/[vp], vp/[np], np/[det, n]]	[det, cat]	4
16.	[s/[vp], vp/[np], np/[n]]	[cat]	3
17.	[s/[vp], vp/[np], np/[n], n/[cat]]	[cat]	
18.	[s/[vp], vp/[np], np/[n]]	[n]	
19.	[s/[vp], vp/[np]]	[np]	
20.	[s/[vp]]	[vp]	
21.	[]	[s]	4

Top down recognizer

1: Push

Looking for Q

Put it on the stack

For convenience, write
 $A \rightarrow B C \dots$
 as
 $A/[B, C \dots]$

```

move([P/[Q | R | Stack], String,           ❶
      [P/[Q | R, Q/S] | Stack], String)
:- rule(Q/S) ← A rule that could find Q
move([P/[Q | R | Stack], [S | String],       ❷
      [P/[Q | R | Stack], [Q | String])
:- rule(Q/[S]).
move([P/[Q | R | S | Stack], [Q | String],   ❸
      [P/[R | S | Stack], String).
move([P/[Q] | Stack, [Q | String],           ❹
      Stack, [P | String]).
  
```

2: Reduce with rule

Looking for Q

```

move([P/[Q | R | Stack], String,           ❶
      [Q/S, P/[Q | R | Stack], String) ← This would satisfy the rule
:- rule(Q/S).
move([P/[Q | R | Stack], [S | String],       ❷
      [P/[Q | R | Stack], [Q | String]) ← So replace it
:- rule(Q/[S]). ← A rule that could find Q
move([P/[Q | R | S | Stack], [Q | String],   ❸
      [P/[R | S | Stack], String).
move([P/[Q] | Stack, [Q | String],           ❹
      Stack, [P | String]).
  
```

3: Pop

Looking for Q

```

move([P/[Q | R] | Stack], String,           ❶
      [Q/S, P/[Q | R] | Stack], String)
:- rule(Q/S).
move([P/[Q | R] | Stack], [S | String],       ❷
      [P/[Q | R] | Stack], [Q | String])
:- rule(Q/[S]).
move([P/[Q | R | S] | Stack], [Q | String],   ❸
      [P/[R | S] | Stack], String).
move([P/[Q] | Stack, [Q | String],           ❹
      Stack, [P | String]). ← Here it is
  
```

So delete both of them

4: Complete Rule

Looking for Q

```

move([P/[Q | R] | Stack], String,           ❶
      [Q/S, P/[Q | R] | Stack], String)
:- rule(Q/S).
move([P/[Q | R] | Stack], [S | String],       ❷
      [P/[Q | R] | Stack], [Q | String])
:- rule(Q/[S]).
move([P/[Q | R | S] | Stack], [Q | String],   ❸
      [P/[R | S] | Stack], String).
move([P/[Q] | Stack, [Q | String],           ❹
      Stack, [P | String]). ← This will satisfy the whole rule
  
```

So delete both of them

This will satisfy the whole rule

The Interpreter

```
parse(Stack0, String0) :-  
    move(Stack0, String0, Stack, String),  
    parse(Stack, String).  
parse([], []).
```

Three Main Paradigms

- Top-down
- Bottom-up
 - Shift-reduce
 - Left-corner

We can manage with less than four cases

Top-down (again)

```
parse([], String, String).  
parse([Goal | Goals], [Goal | String], Rest) :-  
    parse(Goals, String, Rest).  
parse([Goal | Goals], String0, Rest) :-  
    rule(Goal, Rhs),  
    parse(Rhs, String0, String1),  
    parse(Goals, String1, Rest).
```

The empty set of goals is satisfied by the empty string, leaving everything else as the remainder

Top-down

```
parse([], String, String).
parse([Goal | Goals], [Goal | String], Rest) :-
    parse(Goals, String, Rest).
parse([Goal | Goals], String0, Rest) :-
    rule(Goal, Rhs),
    parse(RHS, String0, String1),
    parse(Goals, String1, Rest).
```

If the Goal is found at the beginning of the string, try to show that the remainder of the string satisfies the remainder of the goals.

Top-down

```
parse([], String, String).
parse([Goal | Goals], [Goal | String], Rest) :-
    parse(Goals, String, Rest).
parse([Goal | Goals], String0, Rest) :-
    rule(Goal, Rhs),
    parse(RHS, String0, String1),
    parse(Goals, String1, Rest).
```

If there is a rule that would replace the next goal by some more specific goals (RHS), then

Top-down

```
parse([], String, String).
parse([Goal | Goals], [Goal | String], Rest) :-
    parse(Goals, String, Rest).
parse([Goal | Goals], String0, Rest) :-
    rule(Goal, Rhs),
    parse(RHS, String0, String1),
    parse(Goals, String1, Rest).
```

If there is a rule that would replace the next goal by some more specific goals (RHS), then try to satisfy these more specific goals, and then

Top-down

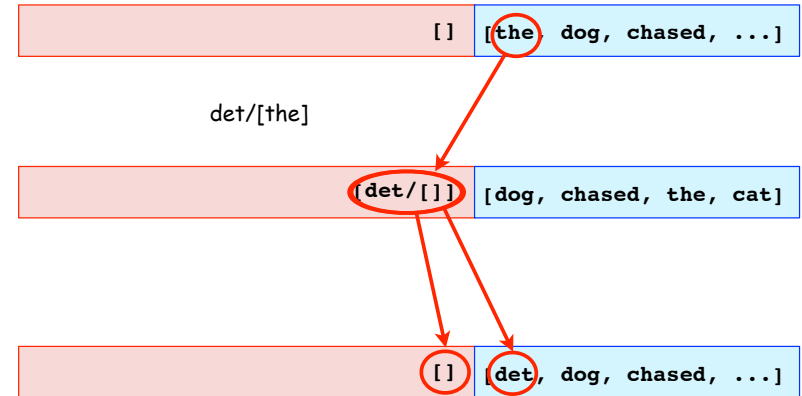
```
parse([], String, String).
parse([Goal | Goals], [Goal | String], Rest) :-
    parse(Goals, String, Rest).
parse([Goal | Goals], String0, Rest) :-
    rule(Goal, Rhs),
    parse(RHS, String0, String1),
    parse(Goals, String1, Rest).
```

If there is a rule that would replace the next goal by some more specific goals (RHS), then try to satisfy these more specific goals, and then try to satisfy the rest of the original set of goals

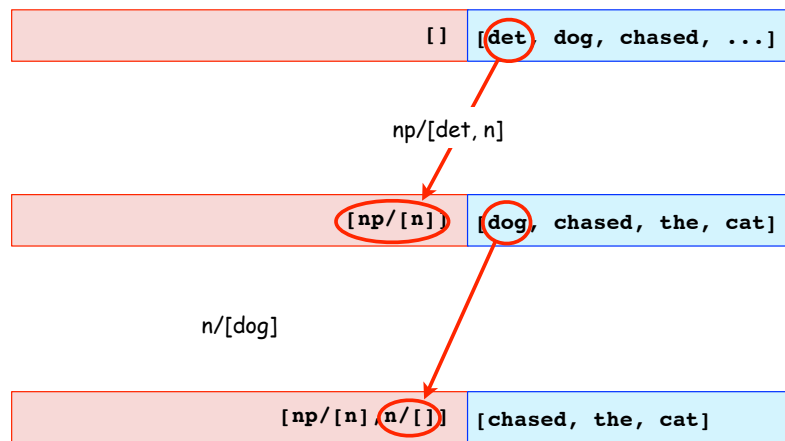
Backtracking

0.		[s/[np, vp]]	[the, dog, chased, the, cat]	0
1.		[np/[det, n], s/[np, vp]]	[the, dog, chased, the, cat]	1
2.		[det/[the], np/[det, n], s/[np, vp]]	[the, dog, chased, the, cat]	1
3.		[np/[det, n], s/[np, vp]]	[det, dog, chased, the, cat]	4
4.		[det/[the], np/[det, n], s/[np, vp]]	[det, dog, chased, the, cat]	1
4.		[np/[n], s/[np, vp]]	[dog, chased, the, cat]	3
5.		[n/[cat], np/[n], s/[np, vp]]	[dog, chased, the, cat]	1
5.		[n/[dog], np/[n], s/[np, vp]]	[dog, chased, the, cat]	1
6.		[np/[n], s/[np, vp]]	[n, chased, the, cat]	4
7.		[n/[cat], np/[n], s/[np, vp]]	[n, chased, the, cat]	1
7.		[n/[dog], np/[n], s/[np, vp]]	[n, chased, the, cat]	1
7.		[n/[house], np/[n], s/[np, vp]]	[n, chased, the, cat]	1
7.		[s/[np, vp]]	[np, chased, the, cat]	4
8.		[np/[det, n], s/[np, vp]]	[np, chased, the, cat]	1
...				

Bottom-up



Bottom-up



0.		[]	[the, dog, chased, the, cat]	0
1.		[det/[the]]	[dog, chased, the, cat]	2
2.		[]	[det, dog, chased, the, cat]	1
3.		[np/[n]]	[dog, chased, the, cat]	2
4.		[np/[n], n/[]]	[chased, the, cat]	2
5.		[np/[n]]	[n, chased, the, cat]	1
6.		[np/[]]	[chased, the, cat]	3
7.		[]	[np, chased, the, cat]	1
8.		[s/[vp]]	[chased, the, cat]	2
9.		[s/[vp], v/[]]	[the, cat]	2
10.		[s/[vp]]	[v, the, cat]	1
11.		[s/[vp], vp/[np]]	[the, cat]	2
12.		[s/[vp], vp/[np], det/[]]	[cat]	2
13.		[s/[vp], vp/[np]]	[det, cat]	1
14.		[s/[vp], vp/[np], np/[n]]	[cat]	2
15.		[s/[vp], vp/[np], np/[n], n/[]]	[]	2
16.		[s/[vp], vp/[np], np/[n]]	[n]	1
17.		[s/[vp], vp/[np], np/[]]	[]	3
18.		[s/[vp], vp/[np]]	[np]	1
19.		[s/[vp], vp/[np], s/[vp]]	[]	2
19.		[s/[vp], vp/[]]	[]	3
20.		[s/[vp]]	[vp]	1
21.		[s/[]]	[]	3
22.		[]	[s]	1

Left Corner

This rule is satisfied

So put the result in the string

```
move(P/[] | Stack], String,
     Stack, P |String]) :- !.
move(Stack, [Q | String],
     [P/R | Stack], String) :- rule(P/[Q | R]).
move([P/[Q | R] | Stack], [Q | String],
     [P/R | Stack], String).
```

Left Corner

Need to consume a Q

Here is a rule that could do it

```
move([P/[] | Stack], String,
     Stack, [P |String]) :- !.
move(Stack, [Q | String],
     [P/R | Stack], String) :- rule(P/[Q | R]).
move([P/[Q | R] | Stack], [Q | String],
     [P/R | Stack], String).
```

No more Q

Try to satisfy rest of the rule

Left Corner

Need to consume a Q

Here is one

```
move([P/[] | Stack], String,
     Stack, [P |String]) :- !.
move(Stack, [Q | String],
     [P/R | Stack], String) :- rule(P/[Q | R]).
move([P/[Q | R] | Stack], [Q | String],
     [P/R | Stack], String).
```

No more Q

Left-corner

Same as for top-down.

```
parse([], String, String).
parse([Goal | Goals], [Goal | String], Rest) :-
  parse(Goals, String, Rest).
parse(Goals, [Item | String0], Rest) :-
  rule(P, [Item | Rhs]),
  parse(RHS, String0, String1),
  parse(Goals, [P | String1], Rest).
```

Left-corner

Find a rule that could form a phrase beginning with the leftmost item in the string (the left corner).

```
parse([], String, String).
parse([Goal | Goals], [Goal | String], Rest) :-
  parse(Goals, String, Rest).
parse(Goals, [Item | String0], Rest) :-
  rule(P, [Item | Rhs]),
  parse(RHS, String0, String1),
  parse(Goals, [P | String1], Rest).
```

Martin Kay

Left-corner Parsing

Left-corner

Find a rule that could form a phrase beginning with the leftmost item in the string (the left corner).

With the right-hand side of the rule as a sequence of goals, try to find such a phrase.

```
parse([], String, String).
parse([Goal | Goals], [Goal | String], Rest) :-
  parse(Goals, String, Rest).
parse(Goals, [Item | String0], Rest) :-
  rule(P, [Item | Rhs]),
  parse(RHS, String0, String1),
  parse(Goals, [P | String1], Rest).
```

Martin Kay

Left-corner Parsing

Left-corner

Find a rule that could form a phrase beginning with the leftmost item in the string (the left corner).

With the right-hand side of the rule as a sequence of goals, try to find such a phrase.

Replace the phrase with its category, and continue trying to satisfy the original goals.

```
parse([], String, String).
parse([Goal | Goals], [Goal | String], Rest) :-
  parse(Goals, String, Rest).
parse(Goals, [Item | String0], Rest) :-
  rule(P, [Item | RHS]),
  parse(RHS, String0, String1),
  parse(Goals, [P | String1], Rest).
```

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Left-corner Parsing

Top-down vs. Left-corner

```
parse([], String, String).
parse([Goal | Goals], [Goal | String], Rest) :-
  parse(Goals, String, Rest).
parse([Goal | Goals], String0, Rest) :-
  rule(Goal, RHS),
  parse(RHS, String0, String1),
  parse(Goals, String1, Rest).
```

```
parse([], String, String).
parse([Goal | Goals], [Goal | String], Rest) :-
  parse(Goals, String, Rest).
parse(Goals, [Item | String0], Rest) :-
  rule(P, [Item | RHS]),
  parse(RHS, String0, String1),
  parse(Goals, [P | String1], Rest).
```

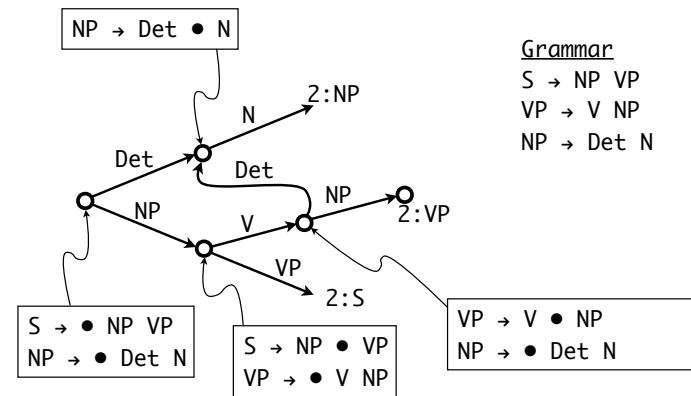
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Left-corner Parsing

Shift reduce

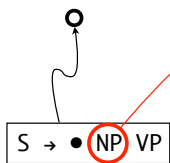
	the dog chased the cat
det	dog chased the cat
det n	chased the cat
np	chased the cat
np v	the cat
np v det	cat
np v det n	
np v np	
np vp	
s	

Finite-state guide



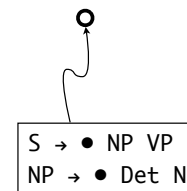
Grammar

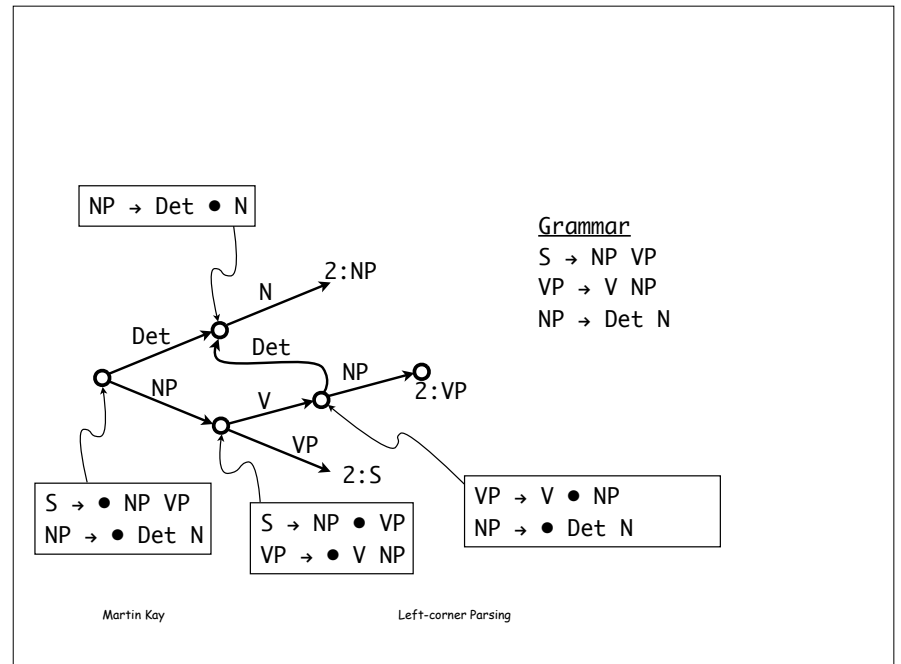
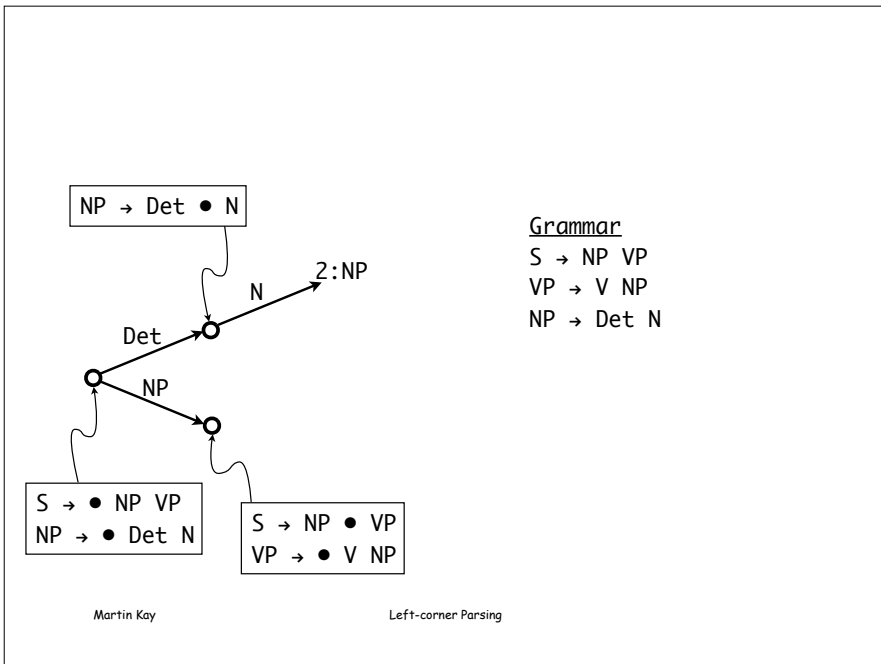
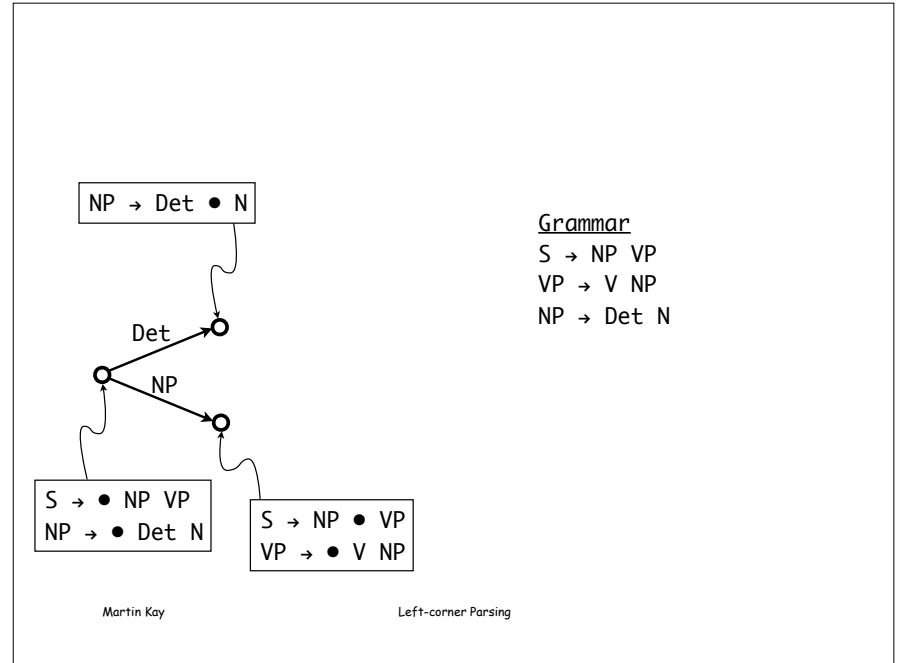
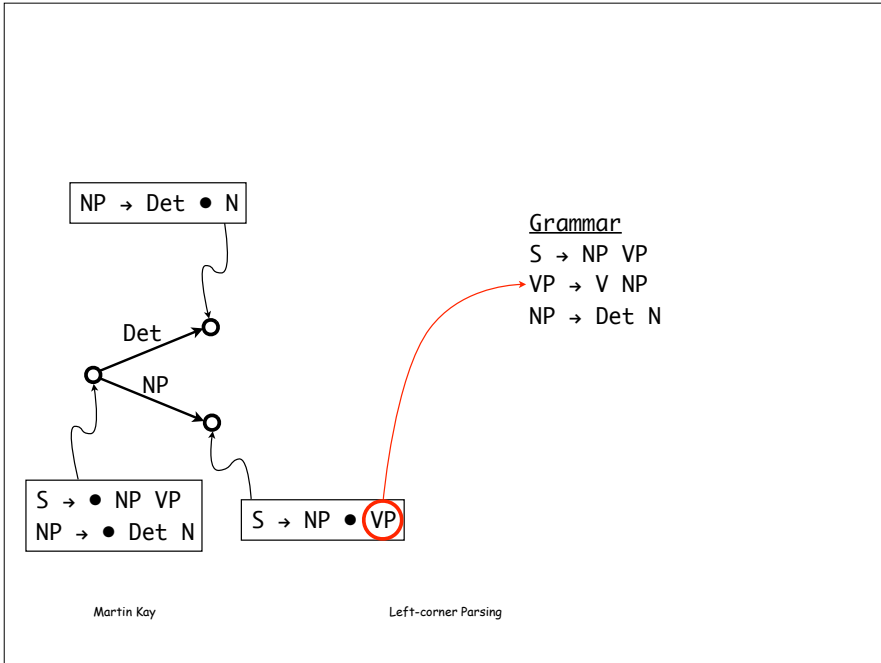
$S \rightarrow NP VP$
 $VP \rightarrow V NP$
 $NP \rightarrow Det N$



Grammar

$S \rightarrow NP VP$
 $VP \rightarrow V NP$
 $NP \rightarrow Det N$





Reachability—Top down

```
parse([], String, String).
parse([Goal | Goals], [Goal | String], Rest) :-
    parse(Goals, String, Rest).
parse([Goal | Goals], String0, Rest) :-
    rule(Goal, Rhs),
    parse(RHS, String0, String1),
    parse(Goals, String1, Rest).
```

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Left-corner Parsing

Reachability—Top down

```
parse([], String, String).
parse([Goal | Goals], [Goal | String], Rest) :-
    parse(Goals, String, Rest).
parse([Goal | Goals], [Item1 | String0], Rest) :-
    rule(Goal, [Item2 | Rhs]),
    reachable(Item2, Item1),
    parse(RHS, [Item1 | String0], String1),
    parse(Goals, String1, Rest).
```

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Left-corner Parsing

Reachability—Bottom-up

```
parse([], String, String).
parse([Goal | Goals], [Goal | String], Rest) :-
    parse(Goals, String, Rest).
parse(Goals, [Item | String0], Rest) :-
    rule(P, [Item | Rhs]),
    parse(RHS, String0, String1),
    parse(Goals, [P | String1], Rest).
```

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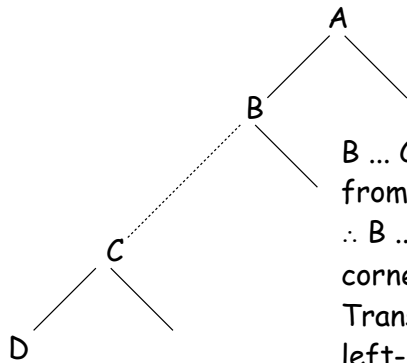
Left-corner Parsing

Reachability—Bottom-up

```
parse([], String, String).
parse([Goal | Goals], [Goal | String], Rest) :-
    parse(Goals, String, Rest).
parse([Goal | Goals], [Item | String0], Rest) :-
    rule(P, [Item | Rhs]),
    reachable(Goal, P),
    parse(RHS, String0, String1),
    parse([Goal | Goals], [P | String1], Rest).
```

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Left-corner Parsing



B ... C, D are reachable from A.
 \therefore B ... C, D are (proper) left corners of A.
 Transitive closure of the left-daughter relation.

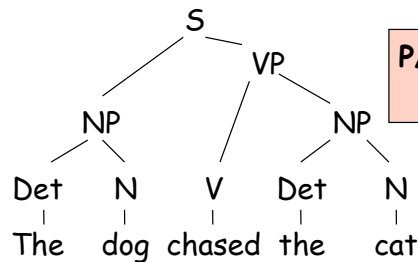
Remember this?

A grammar that contains instances of left recursion can be replaced by one with no left recursion which

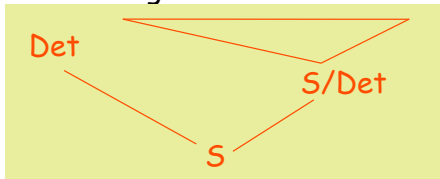
- describes the same language
- assigns the same structure to the sentences.

COMING SOON!

The Left-corner Transformation



P/Q: The part of a P that follows a Q



Rules

```
s --> np, vp.
np --> det, n.
n --> adj, n.
np --> np, pp.
det --> np, ap_s.
pp --> prep, np.
vp --> iv.
vp --> vt, np.
vp --> vp, pp.
```

```
det --> [the].
n --> [dog].
n --> [cat]
n --> [bone].
n --> [field].
adj --> [big].
prep --> [with].
prep --> [in].
vi --> [slept].
vt --> [chased].
```

Eliminating Left Recursion (step 1)

Reachability

```
s:   np
      det
np:  det
det: det
n:   n
vp:  iv
      tv
pp:  prep
```

```
s --> np, s/np.
s --> det, s/det.
np --> det, np/det.
np --> np, np/np.
det --> np, det/np.
det --> det, det/det.
pp --> prep, pp/prep.
vp --> iv, vp/iv.
vp --> vt, vp/vt.
vp --> vp, vp/vp
```

Rules

```

s --> np, vp.
np --> det, n.
n --> adj, n.
np --> np, pp.
det --> np, ap_s.
pp --> prep, np.
vp --> iv.
vp --> vt, np.
vp --> vp, pp.

det --> [the].
n --> [dog].
n --> [cat].
n --> [bone].
n --> [field].
adj --> [big].
prep --> [with].
prep --> [in].
vi --> [slept].
vt --> [chased].

```

Step 2

```

s --> np, s/np.
s --> det, s/det.
np --> det, np/det.
np --> np, np/np.
det --> np, det/np.
det --> det, det/det.
pp --> prep, pp/prep.
vp --> iv, vp/iv.
vp --> vt, vp/vt.
vp --> vp, vp/vp

```

```

s/np --> vp, s/s.
s/det --> n, s/np.
np/det --> n, np/np.
np/np --> [].
np/np --> ap/s, np/det.
det/np --> ap_s, det/det.
pp/prep --> np, pp/pp
...

```

Martin Kay Left-corner Parsing

Rules

```

s --> np, vp.
np --> det, n.
n --> adj, n.
np --> np, pp.
det --> np, ap_s.
pp --> prep, np.
vp --> iv.
vp --> vt, np.
vp --> vp, pp.

det --> [the].
n --> [dog].
n --> [cat].
n --> [bone].
n --> [field].
adj --> [big].
prep --> [with].
prep --> [in].
vi --> [slept].
vt --> [chased].

```

Step 2

```

s --> np, s/np.
s --> det, s/det.
np --> det, np/det.
np --> np, np/np.
det --> np, det/np.
det --> det, det/det.
pp --> prep, pp/prep.
vp --> iv, vp/iv.
vp --> vt, vp/vt.
vp --> vp, vp/vp

```

```

s/np --> vp, s/s.
s/det --> n, s/np.
np/det --> n, np/np.
np/np --> [].
np/np --> ap_s, np/det.
det/np --> ap_s, det/det.
pp/prep --> np, pp/pp
...

```

Martin Kay Left-corner Parsing

Caching

```

foo(Input1, Input2, Output1, Output2)

```

Inputs Outputs

```

try(foo(A, B, C, F)) :-
  \+ done(foo(A, B)),
  assert(done(foo(A, B))),
  foo(A, B, C, D),
  assert(result(foo(A, B, C, D))),
  fail.
try(foo(A, B, C, D)) :-
  result(foo(A, B, C, D)).

```

Find all solutions on the first call!

Martin Kay Left-corner Parsing

Caching top-down

```

parse([], String, String).
parser(Goals, String, String) :-
  retract(parse(Goals, String, String)),
  retract(parse(Goals, String, String)),
  try(parse(Goals, String, String)) :-
    rule(Goal, RHS),
    try(parse(RHS, String0, String1)),
    try(parse(Goals, String1, Rest)).
try(parse(Goals, String0, String)) :-
  \+ done(parse(Goals, String0)),
  assert(done(parse(Goals, String0))),
  parse(Goals, String0, String),
  assert(result(parse(Goals, String0, String))),
  fail.
try(parse(Goals, String0, String)) :-
  result(parse(Goals, String0, String)).

```

Martin Kay Left-corner Parsing

A repackaged left-corner parser

```
parse([], String, String).
parse([Goal | Goals], [Goal | String0], String) :-
    parse(Goals, String0, String).
parse(Goals, String0, String) :-
    apply_rule(String0, String1),
    parse(Goals, String1, String).

apply_rule([Item | String0], [Phrase | String]) :-
    rule(Phrase, [Item | Rhs]),
    parse(Rhs, String0, String),
```

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Left-corner Parsing

apply_rule as a memo-function

```
apply_rule(String0, String) :-
    table(String0, Strings),
    !,
    member(String, Strings).
apply_rule(String0, String) :-
    setof(S, apply_rule0(String0, S), Strings),
    assert(table(String0, Strings)),
    member(String, Strings).

apply_rule0([Item | String0], [Phrase | String]) :-
    rule(Phrase, [Item | Rhs]),
    parse(Rhs, String0, String).
```

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Left-corner Parsing

The Cache

```
table([the,dog,chased,the,cat],
      [det,dog,chased,the,cat]).
table([dog,chased,the,cat], [n,chased,the,cat]).
table([det,dog,chased,the,cat],
      [np,chased,the,cat]).
table([chased,the,cat], [v,the,cat]).
table([the,cat], [det,cat]).
table([cat], [n]).
table([det,cat], [np]).
table([v,the,cat], [vp]).
table([np,chased,the,cat], [s]).
```

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Left-corner Parsing

- **Bottom-up parsers also use a stack, but in this case, the stack represents a summary of the input already seen, rather than a prediction about input yet to be seen.**

Martin Kay

Left-corner Parsing

Backtracking Complexity

- Exponential—why?