

Phase Transitions in Semidefinite Relaxations

Andrea Montanari

[with Adel Javanmard, Federico Ricci-Tersenghi, Subhabrata Sen]

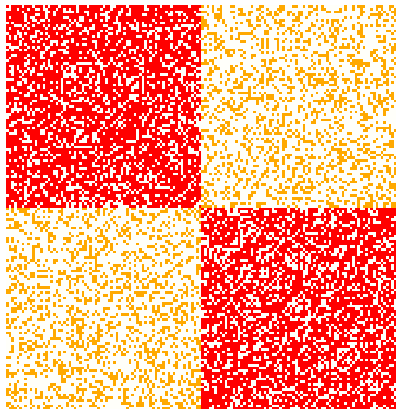
Stanford University

April 5, 2016

What is this talk about?

SDP for Matrix/Graph estimation

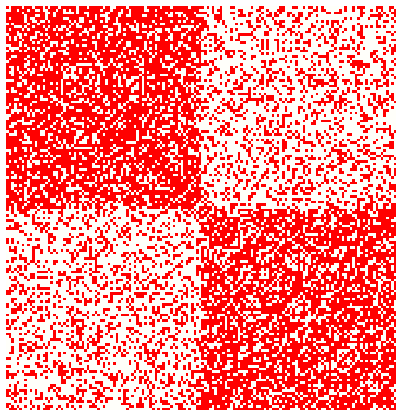
The hidden partition model



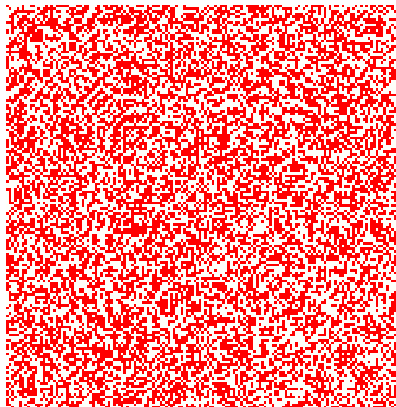
Vertices V , $|V| = n$, $V = V_+ \cup V_-$, $|V_+| = |V_-| = n/2$

$$\mathbb{P}\{(i, j) \in E\} = \begin{cases} p & \text{if } \{i, j\} \subseteq V_+ \text{ or } \{i, j\} \subseteq V_-, \\ q < p & \text{otherwise.} \end{cases}$$

Of course entries are not colored. . .



...and rows/columns are not ordered



Problem: Detect/estimate the partition

What is this talk about?

SDP for Matrix/Graph estimation

Exact phase transition(?)

Outline

[arXiv:1603.1603.04064]
[arXiv:1511.08769, PNAS]
[arXiv:1504.05910, STOC]

Background

Estimation

$$\mathbf{x}_0 \in \{+1, -1\}^n, \langle \mathbf{x}_0, \mathbf{1} \rangle = 0$$

$$\mathbb{P}\{(i, j) \in E\} = \begin{cases} p & \text{if } x_{0,i} = x_{0,j}, \\ q & \text{otherwise.} \end{cases}$$

Estimation ($p = a/n, q = b/n$)

$$\mathbf{x}_0 \in \{+1, -1\}^n, \langle \mathbf{x}_0, \mathbf{1} \rangle = 0$$

$$\mathbb{P}\{(i, j) \in E\} = \begin{cases} a/n & \text{if } x_{0,i} = x_{0,j}, \\ b/n & \text{otherwise.} \end{cases}$$

Estimation threshold

$$\mathbf{x}_0 \in \{+1, -1\}^n, \langle \mathbf{x}_0, \mathbf{1} \rangle = 0$$

$$\mathbb{P}\{(i, j) \in E\} = \begin{cases} a/n & \text{if } x_{0,i} = x_{0,j}, \\ b/n & \text{otherwise.} \end{cases}$$

$$\hat{\mathbf{x}} : \mathcal{G} \rightarrow \{+1, -1\}^n,$$

$$\text{Overlap}_n(\hat{\mathbf{x}}) = \frac{1}{n} \mathbb{E}\{|\langle \hat{\mathbf{x}}(G), \mathbf{x}_0 \rangle|\}.$$

For which a, b , $\liminf_{n \rightarrow \infty} \text{Overlap}_n(\hat{\mathbf{x}}) \geq \delta > 0$?

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For which a, b , $\liminf_{n \rightarrow \infty} \text{Overlap}_n(\hat{\mathbf{x}}) \geq \delta > 0$?

Information theory threshold

Theorem (Mossel, Neeman, Sly, 2012)

There is an estimator that is better than random if and only if

$$\frac{a - b}{\sqrt{2(a + b)}} > 1.$$

[Proves conjecture by Decelle, Krzakala, Moore, Zdeborova, 2011]

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Computational threshold

- ▶ Dyer, Frieze 1989
- ▶ Condon, Karp 2001
- ▶ McSherry 2001
- ▶ Coja-Oghlan 2010
- ▶ Massoulié 2013 and Mossel, Neeman, Sly, 2013

$$p = na > q = nb \text{ fixed.}$$

$$a - b \gg n^{1/2}$$

$$a - b \gg \sqrt{b \log n}$$

$$a - b \gg \sqrt{b}$$

$$\frac{a - b}{\sqrt{2(a + b)}} > 1$$

Very ingenious spectral methods!

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Very ingenious spectral methods!

What if I am not ingenious?

Maximum Likelihood

Adjacency matrix

$$A_{ij} = \begin{cases} 1 & \text{if } (i, j) \in E, \\ 0 & \text{otherwise.} \end{cases}$$

$$\mathbf{A} = (A_{ij})_{1 \leq i, j \leq n}$$

Maximum likelihood

$$\sigma_i = \begin{cases} +1 & \text{if } i \in V_+, \\ -1 & \text{if } i \in V_-. \end{cases}$$

$$\begin{aligned} &\text{maximize} && \sum_{i,j=1}^n A_{ij} \sigma_i \sigma_j, \\ &\text{subject to} && \sum_{i=1}^n \sigma_i = 0, \\ &&& \sigma_i \in \{+1, -1\}. \end{aligned}$$

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Centered adjacency matrix

$$A_{ij}^{\text{cen}} = \begin{cases} 1 - (d/n) & \text{if } (i, j) \in E, \\ -(d/n) & \text{otherwise.} \end{cases}$$

$$A^{\text{cen}} = A - \frac{d}{n} \mathbf{1} \mathbf{1}^T$$

Lagrangian

$$\begin{aligned} & \text{maximize} && \langle \mathbf{A}^{\text{cen}}, \boldsymbol{\sigma} \boldsymbol{\sigma}^T \rangle, \\ & \text{subject to} && \boldsymbol{\sigma} \in \{+1, -1\}^n. \end{aligned}$$

- ▶ NP-hard
- ▶ $\text{SDP}(\mathbf{A}^{\text{cen}})$ is a very natural convex relaxation

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- ▶ NP-hard
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Relaxation

$$\begin{aligned} & \text{maximize} && \langle \mathbf{A}^{\text{cen}}, \boldsymbol{\sigma} \boldsymbol{\sigma}^T \rangle, \\ & \text{subject to} && \boldsymbol{\sigma} \in \{+1, -1\}^n. \end{aligned}$$

SDP(\mathbf{A}^{cen}):

$$\begin{aligned} & \text{maximize} && \langle \mathbf{A}^{\text{cen}}, \mathbf{X} \rangle, \\ & \text{subject to} && \mathbf{X} \in \mathbb{R}^{n \times n}, \mathbf{X} \succeq 0, \\ & && X_{ii} = 1. \end{aligned}$$

- ▶ This is really **off-the-shelf**
- ▶ How well does it work?

Near-optimality of SDP

Before we pass to SDP

- ▶ What's the problem with sparse graphs?
- ▶ What's the problem vanilla PCA?

Why PCA?

Ground truth

$$x_{0,i} = \begin{cases} +1 & \text{if } i \in V_+, \\ -1 & \text{if } i \in V_-. \end{cases}$$

Data = RankOne + Wigner

$$\frac{1}{\sqrt{d}} A^{\text{cen}} = \frac{\lambda}{n} x_0 x_0^\top + W, \quad \lambda \equiv \frac{a - b}{\sqrt{2(a + b)}}$$

$$\mathbb{E}\{W_{ij}\} = 0, \quad \mathbb{E}\{W_{ij}^2\} \in \left\{ \frac{a}{dn}, \frac{b}{dn} \right\} \approx \frac{1}{n}.$$

Why PCA?

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The *right* parametrization

$$d = \frac{a + b}{2}, \quad \lambda = \frac{a - b}{\sqrt{2(a + b)}}$$

Naive PCA

$$\hat{\mathbf{x}}^{\text{PCA}}(\mathbf{A}^{\text{cen}}) = \text{sign}(\mathbf{v}_1(\mathbf{A}^{\text{cen}})).$$

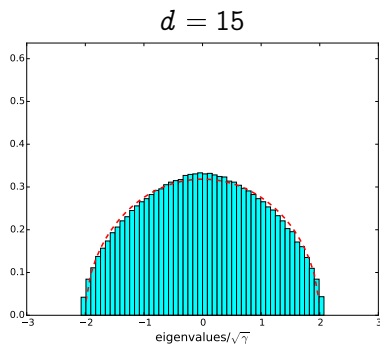
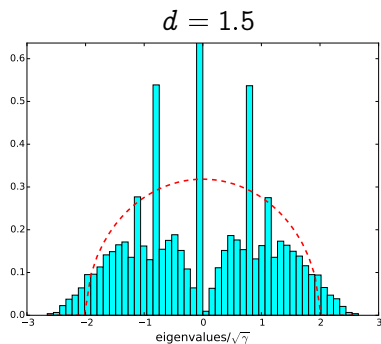
Does it work?

No!

Does it work?

No!

Spectral relaxation bad in the sparse regime!

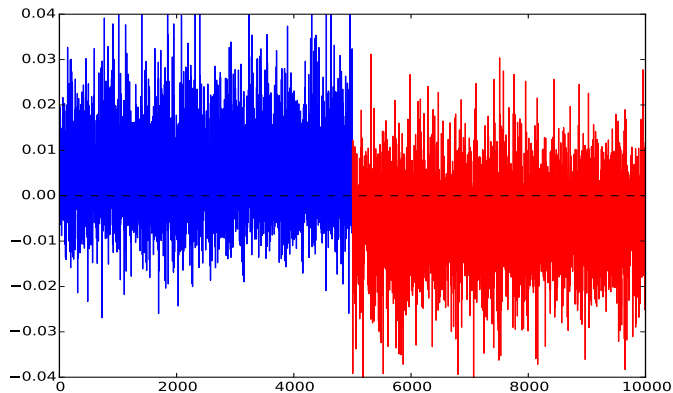


Theorem (Krivelevich, Sudakov 2003)

With high probability, for $d = O(1)$, $\lambda = 0$

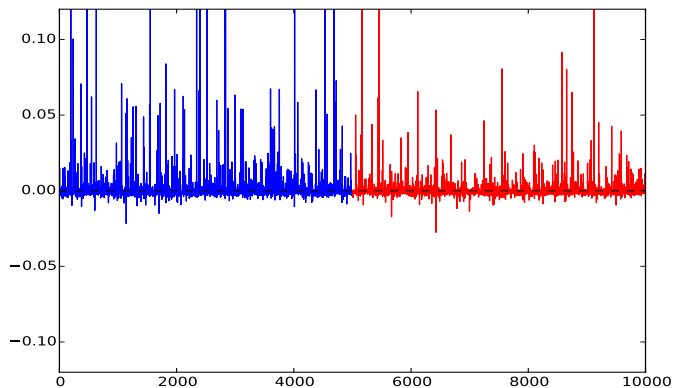
$$\lambda_{\max}(\mathbf{A}^{cen}/\sqrt{d}) = C\sqrt{\log n/(\log \log n)}(1 + o(1)).$$

Example: $d = 20$, $\lambda = 1.2$, $n = 10^4$



$v_1(A^{\text{cen}})$

Example: $d = 3$, $\lambda = 1.2$, $n = 10^4$



$v_1(A^{\text{cen}})$

Why should SDP work better?

$$\begin{aligned} & \text{maximize} && \langle \mathbf{A}^{\text{cen}}, \mathbf{X} \rangle, \\ & \text{subject to} && \mathbf{X} \in \mathbb{R}^{n \times n}, \mathbf{X} \succeq \mathbf{0}, \\ & && X_{ii} = 1. \end{aligned}$$

Recall the info-theoretic limit

$$d = \frac{a + b}{2} > 1, \quad \lambda = \frac{a - b}{\sqrt{2(a + b)}}$$

Theorem (Mossel, Neeman, Sly, 2012)

There is an estimator that is better than random if and only if

$$\lambda > 1.$$

SDP has nearly optimal threshold

Theorem (Montanari, Sen 2015)

The SDP estimator is better than random if^a

$$\lambda \geq 1 + o_d(1).$$

^aCan replace $o_d(1)$ by $C_0 d^{-c_1}$.

SDP estimator

Ideally

- ▶ Let $\mathbf{X}_* \in \mathbb{R}^{n \times n}$ be an optimizer
- ▶ Compute the principal eigen-vector $\mathbf{v}_1(\mathbf{X}_*)$.
- ▶ Return $\hat{\mathbf{x}}^{\text{SDP}}(G) = \text{sign}(\mathbf{v}_1(\mathbf{X}_*))$

In reality, a few tricks

SDP estimator

Ideally

- ▶ Let $X_* \in \mathbb{R}^{n \times n}$ be an optimizer
- ▶ Compute the principal eigen-vector $v_1(X_*)$.
- ▶ Return $\hat{x}^{\text{SDP}}(G) = \text{sign}(v_1(X_*))$

In reality, a few tricks

Key technical result: SDP value

Theorem (Montanari, Sen 2015)

Assume $G \sim G(n, d, \lambda)$.

If $\lambda \leq 1$, then, with high probability,

$$\frac{1}{n\sqrt{d}} \text{SDP}(A_G^{\text{cen}}) = 2 + o_d(1).$$

If $\lambda > 1$, then there exists $\Delta(\lambda) > 0$ such that, with high probability,

$$\frac{1}{n\sqrt{d}} \text{SDP}(A_G^{\text{cen}}) = 2 + \Delta(\lambda) + o_d(1).$$

Earlier/related work

Optimal spectral methods

- ▶ Massoulié 2013
- ▶ Mossel, Neeman, Sly, 2013
- ▶ Bordenave, Lelarge, Massoulié, 2015

SDP, $d = \Theta(\log n)$

- ▶ Abbe, Bandeira, Hall 2014
- ▶ Hajek, Wu, Xu 2015

SDP, detection

- ▶ Guédon, Vershynin, 2015 (requires $\lambda \geq 10^4$, very different proof)

How does SDP work 'in practice'?

Thresholds

- ▶ $\lambda_c^{\text{opt}}(d) \equiv$ Threshold for optimal estimator
- ▶ $\lambda_c^{\text{SDP}}(d) \equiv$ Threshold for SDP-based estimator

What we know

▶ $\lambda_c^{\text{opt}}(d) = 1$

[Mossel, Neeman, Sly, 2013]

▶ $\lambda_c^{\text{SDP}}(d) = 1 + o_d(1)$

[Montanari, Sen, 2015]

How big is the $o_d(1)$ gap?

What we know

▶ $\lambda_c^{\text{opt}}(d) = 1$

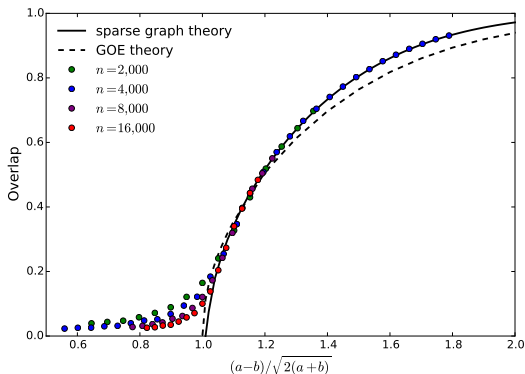
[Mossel, Neeman, Sly, 2013]

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[Montanari, Sen, 2015]

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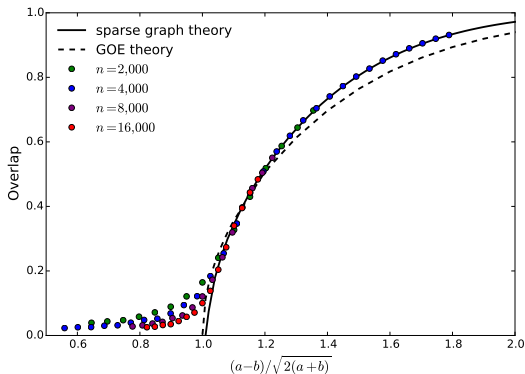
Simulations: $d = 5$, $N_{\text{sample}} = 500$ (with Javanmard and Ricci)



SDP estimator $\hat{\mathbf{x}}^{\text{SDP}} \in \{+1, -1\}^n$

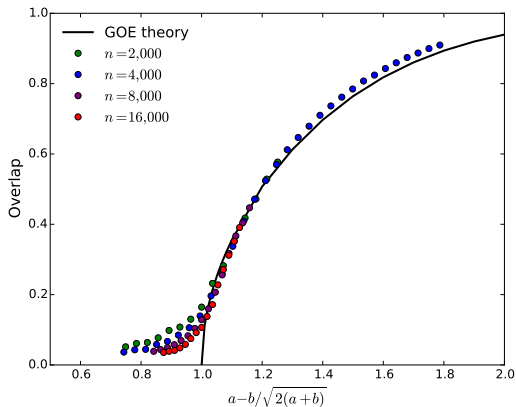
$$\text{Overlap}_n(\hat{\mathbf{x}}) = \frac{1}{n} \mathbb{E} \{ |\langle \hat{\mathbf{x}}^{\text{SDP}}(G), \mathbf{x}_0 \rangle| \}.$$

Simulations: $d = 5$, $N_{\text{sample}} = 500$



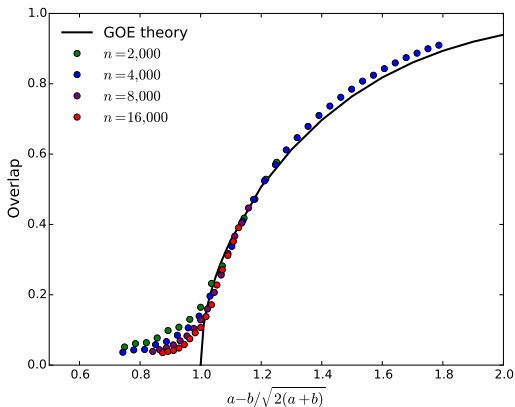
$$\lambda_c^{\text{SDP}}(d = 5) \approx 1.$$

Simulations: $d = 10$, $N_{\text{sample}} = 500$



$$\lambda_c^{\text{SDP}}(d=10) \approx 1.$$

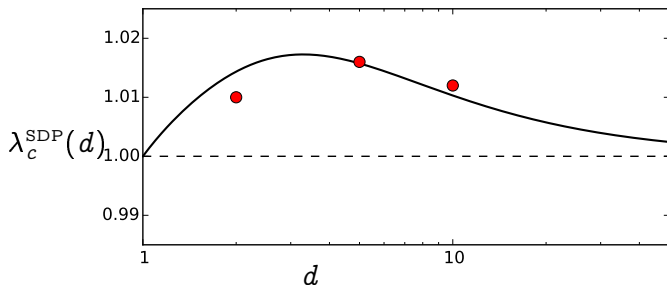
Simulations: $d = 10$, $N_{\text{sample}} = 500$



Estimate $\lambda_c^{SDP}(d)$ from data...

(a long story)

$$\lambda_c^{\text{SDP}}(d)$$



- ▶ Dots: Numerical estimates
- ▶ Line: Non-rigorous analytical approximation (using statistical physics)
- ▶ **At most 2% sub-optimal!**
- ▶ (Moitra, Perry, Wein 2016) suggests $\lambda_c^{\text{SDP}}(d) > 1$ for d small
- ▶ (Fan, Montanari, 2016) $\lambda_c^{\text{SDP}}(d) < c(d) \leq 2$ for all d

How do you actually solve the SDP?

How do you *actually* solve the SDP?

- ▶ Can do $n = 10^5$ in minutes.
- ▶ Code: <http://web.stanford.edu/~montanar/SDPgraph>
- ▶ Burer-Monteiro trick: solve rank-constrained problem

Rank-constrained SDP (non-convex!!!)

$$\begin{aligned} & \text{maximize} && \langle M, X \rangle, \\ & \text{subject to} && X \in \mathbb{R}^{n \times n}, X \succeq 0, \\ & && X_{ii} = 1, \\ & && \text{rank}(X) \leq k \end{aligned}$$

Equivalently:

$$\begin{aligned} & \text{maximize} && F_M(\sigma) \equiv \sum_{i,j=1}^n M_{ij} \langle \sigma_i, \sigma_j \rangle, \\ & \text{subject to} && \sigma_i \in \mathbb{R}^k, \|\sigma_i\|_2 = 1. \end{aligned}$$

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A Grothendieck-type inequality for local maxima

Theorem (Montanari, 2016)

Let σ be a local maximum of F_M over $(S^{k-1})^n$. Then

$$F_M(\sigma) \geq \text{SDP}(M) - \frac{8}{\sqrt{k}} n \|M\|_2.$$

Typically:

- ▶ $\|M\|_2 = O(1)$, $\text{SDP}(M) = \Theta(n)$.
- ▶ **Sufficient to take $k = O(1)$!!!!**

Related work

- ▶ Bandeira, Boumal, Voroninski, 2016:
 \mathbb{Z}_2 synchronization, $k = 2, \lambda > C$

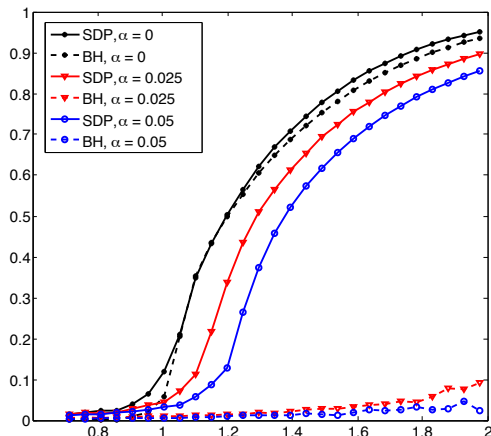
One last question

Is this approach robust to model miss-specifications?

An experiment

- ▶ Select $S \subseteq V$ uniformly at random. with $|S| = n\alpha$.
- ▶ For each $i \in S$, connect all of its neighbors.

An experiment



- ▶ Solid line:SDP
- ▶ Dashed line: Spectral
(Non-backtracking walk [Krzakala, Moore, Mossel, Neeman, Sly, Zdeborova, Zhang, 2013])

Proof ideas

What we want to prove

Theorem (Montanari, Sen 2015)

If $\lambda \leq 1$, then, with high probability,

$$\frac{1}{n\sqrt{d}} \text{SDP}(\mathbf{A}_G^{\text{cen}}) = 2 + o_d(1).$$

If $\lambda > 1$, then there exists $\Delta(\lambda) > 0$ such that, with high probability,

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Strategy

I. Prove equivalence to Gaussian model

II. Analyze Gaussian model

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II. Analyze Gaussian model

Gaussian model: \mathbb{Z}_2 synchronization

$$B(\lambda) \equiv \frac{\lambda}{n} \mathbf{x}_0 \mathbf{x}_0^\top + \mathbf{W}.$$

$\mathbf{x}_0 \in \{+1, -1\}^n$, $\mathbf{W} \sim \text{GOE}(n)$

▶ $(W_{ij})_{i < j} \sim_{iid} \text{N}(0, 1/n)$

▶ $\mathbf{W} = \mathbf{W}^\top$

▶ *A lot is known about spectral properties of B*

Need to characterize the SDP value with Gaussian data

Gaussian model: \mathbb{Z}_2 synchronization

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Need to characterize the SDP value with Gaussian data

Notation

$$\bar{s}(\lambda) \equiv \limsup_{n \rightarrow \infty} \frac{1}{n} \text{SDP}(B(\lambda)),$$

$$\underline{s}(\lambda) \equiv \liminf_{n \rightarrow \infty} \frac{1}{n} \text{SDP}(B(\lambda)).$$

$$\text{SDP}(B) \equiv \max \{ \langle B, X \rangle : X \succeq 0, X_{ii} = 1 \forall i \}$$

Phase transition at $\lambda = 1$!

Theorem (Montanari, Sen, 2015)

The following holds almost surely

$$\lambda \in [0, 1] \Rightarrow \underline{s}(\lambda) = \bar{s}(\lambda) = 2,$$

$$\lambda \in (1, \infty) \Rightarrow \underline{s}(\lambda) > 2 \quad (\textit{strictly}).$$

For explicit probability bounds, see the paper

Proof of Gaussian phase transition

Simple facts:



$$\bar{s}(\lambda) \leq \lim_{n \rightarrow \infty} \sigma_{\max}(B(\lambda)) = \begin{cases} 2 & \text{if } \lambda \in [0, 1], \\ \lambda + \lambda^{-1} & \text{if } \lambda \in (1, \infty). \end{cases}$$

[Baik, Ben Arous, Peche, 2005]



$$\underline{s}(\lambda) \geq \lim_{n \rightarrow \infty} \frac{1}{n} \langle \mathbf{1}, B(\lambda) \mathbf{1} \rangle = \lambda$$

- ▶ $\underline{s}(\lambda), \bar{s}(\lambda)$ are non-random, non-decreasing

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Proof of Gaussian phase transition

Simple facts:



$$\bar{s}(\lambda) \leq \lim_{n \rightarrow \infty} \sigma_{\max}(\mathbf{B}(\lambda)) = \begin{cases} 2 & \text{if } \lambda \in [0, 1], \\ \lambda + \lambda^{-1} & \text{if } \lambda \in (1, \infty). \end{cases}$$

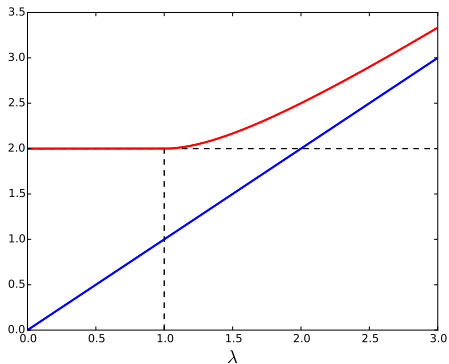
[Baik, Ben Arous, Peche, 2005]



$$\underline{s}(\lambda) \geq \lim_{n \rightarrow \infty} \frac{1}{n} \langle \mathbf{1}, \mathbf{B}(\lambda) \mathbf{1} \rangle = \lambda$$

- ▶ $\underline{s}(\lambda), \bar{s}(\lambda)$ are non-random, non-decreasing

Summarizing



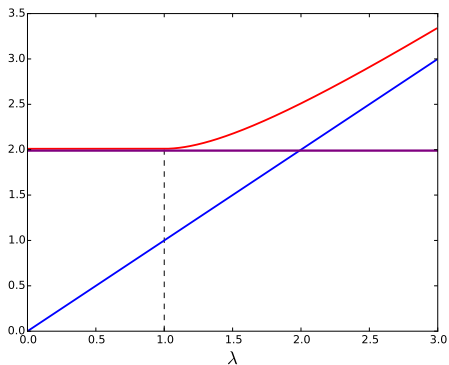
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- ▶ Blue: Lower bound

Proof of Gaussian phase transition

Part 1:

Prove that $\underline{s}(\lambda = 0) \geq 2$

Hence



- ▶ Red: Upper bound
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- ▶ Purple: Non-trivial lower bound

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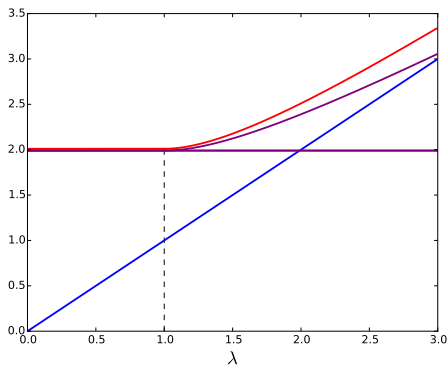
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Prove that $\underline{s}(\lambda = 1 + \varepsilon) > 2$

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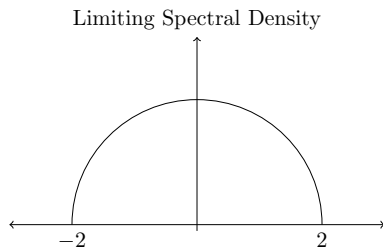
Prove that $\underline{s}(\lambda = 0) \geq 2$

Part 2:

Prove that $\underline{s}(\lambda = 1 + \varepsilon) > 2$

Technique: Construct feasible \mathbf{X} , such that $\langle \mathbf{A}, \mathbf{X} \rangle \geq \dots$

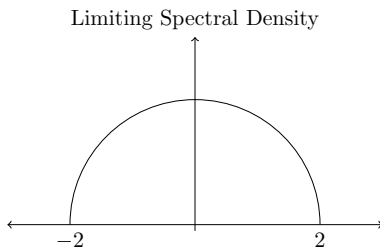
Part 1: $\lambda = 0$



First idea:

- ▶ $v_1 = v_1(B) \equiv$ principal eigenvector of B
- ▶ Take $X = n v_1 v_1^T$
- ▶ Wrong: $X_{ii} \approx N(0, 1)^2 \neq 1$

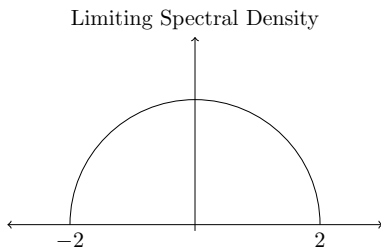
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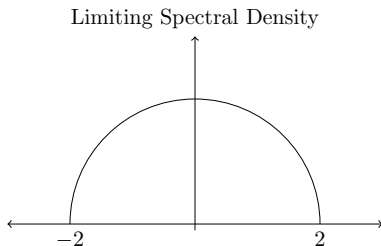
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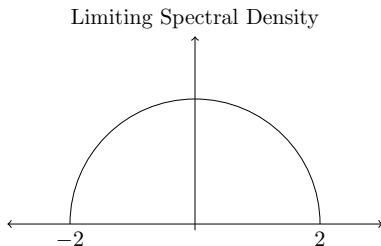
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Good idea:

- ▶ Let $U = [v_1 | v_2 | \cdots | v_{n\delta}] \in \mathbb{R}^{n \times n\delta}$, δ small.
- ▶ $D \equiv \text{Diag}(U U^\top) \in \mathbb{R}^{n \times n}$. Claim $D \approx \delta I$ ($D_{ii} \sim n^{-1} \chi_{n\delta}$)
- ▶ Set $X = D^{-1/2}(U U^\top)D^{-1/2}$

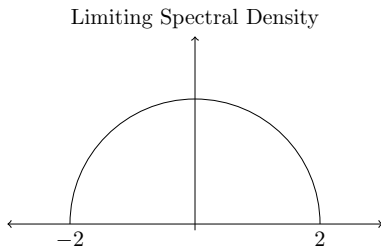
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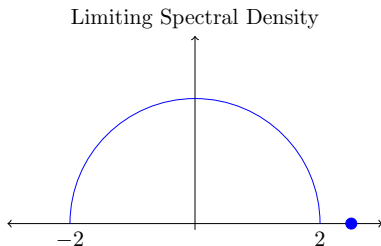
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Part 2: $\lambda = 1 + \varepsilon$



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▶ $T(x) \equiv \max(\min(x, +1), -1), \varphi \in \mathbb{R}^n$

$$\varphi_i \equiv T(\varepsilon\sqrt{n} v_{1,i}).$$

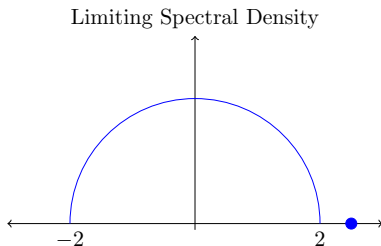
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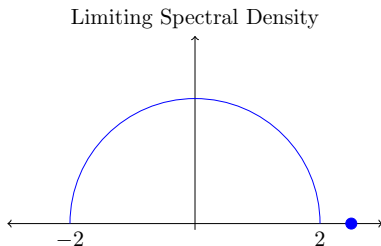
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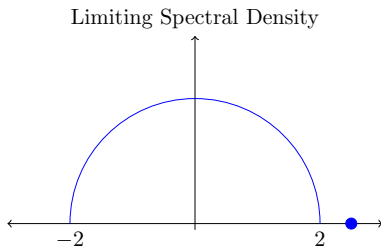
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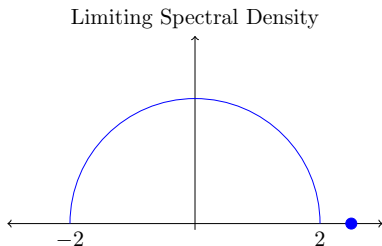


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Strategy

I. Prove equivalence to Gaussian model

II. Analyze Gaussian model

We want to prove

$$\frac{1}{\sqrt{d}} \text{SDP}(\mathbf{A}^{\text{cen}}) \approx \text{SDP}(\mathbf{B}(\lambda))$$

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A simple Lindeberg lemma

- ▶ X_1, X_2, \dots, X_M iid

$$X_i = \begin{cases} \frac{1}{\sqrt{d}} \left(1 - \frac{d}{n}\right) & \text{with probability } \frac{d}{n}, \\ -\frac{\sqrt{d}}{n} & \text{with probability } \left(1 - \frac{d}{n}\right). \end{cases}$$

$$\mathbb{E}\{X_i\} = 0, \mathbb{E}\{X_i^2\} = (1/n) - (d/n^2)$$

- ▶ $Z_1, Z_2, \dots, Z_M \sim_{i.i.d.} \mathcal{N}(0, 1/n)$

Lemma

Assume $M = n(n-1)/2$, $F \in C_3(\mathbb{R}^M)$, $d \leq n^{2/3}/10$. Then

$$\left| \mathbb{E}F(X) - \mathbb{E}F(Z) \right| \leq \frac{n}{3\sqrt{d}} \max_{i \in [M]} \left(\|\partial_i^2 F\|_\infty \vee \|\partial_i^3 F\|_\infty \right).$$

where $\partial_i^\ell F(x) \equiv \frac{\partial^\ell F}{\partial x_i^\ell}$, and $\|\partial_i^\ell F\|_\infty \equiv \sup_{x \in \mathbb{R}^M} |\partial_i^\ell F(x)|$.

Problem: $F(\cdot) = \text{SDP}(\cdot) \notin C_3(\mathbb{R}^M)$

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SDP(A):

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Free energy

$$\Phi_k(\beta, k; A) \equiv \frac{1}{\beta} \log \left\{ \int \exp \left(\beta \sum_{i,j=1}^n A_{i,j} \langle \sigma_i, \sigma_j \rangle \right) \nu_{0,k}(\mathrm{d}\sigma) \right\}$$

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- ▶ Sharp information about eigenvalues of random matrices
- ▶ A lot of work on SDP with random data
[Srebro, Fazel, Parrillo, Candés, Recht, Gross, myself, ...]
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