

Computational Barriers in Statistical Estimation

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Statistical estimation/Statistical learning

Class of models ($\Theta \subseteq \mathbb{R}^d$)

$$\mathcal{C}_\Theta \equiv \{ \mathbb{P}_{\boldsymbol{\theta}} : \quad \boldsymbol{\theta} \in \Theta \}$$

Data

$$x_1, x_2, \dots, x_n \sim_{iid} \mathbb{P}_{\boldsymbol{\theta}_0}(\cdot)$$

Estimate $\boldsymbol{\theta}_0$ from data $x_1^n = (x_1, \dots, x_n)$

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Minimax theory

Loss function

$$L : \mathbb{R}^p \times \mathbb{R}^p \rightarrow \mathbb{R}$$
$$(\hat{\theta}, \theta_0) \mapsto L(\hat{\theta}, \theta_0)$$

Minimax risk

$$R_n^*(\Theta) = \inf_{\hat{\theta}(\cdot)} \sup_{\theta_0 \in \Theta} \mathbb{E}_{\theta_0} L(\hat{\theta}(x_1^n), \theta_0)$$

[Wald, 1950]

Kindergarten example

$$\mathbb{P}_{\boldsymbol{\theta}}(\cdot) = \mathcal{N}(\boldsymbol{\theta}, \mathbf{I}_d)$$

$$L(\boldsymbol{\theta}_0, \hat{\boldsymbol{\theta}}) = \|\boldsymbol{\theta}_0 - \hat{\boldsymbol{\theta}}\|_2^2$$

$$\Theta = \mathbb{R}^d.$$

Theorem

$$R_n^*(\Theta) = \frac{d}{n}.$$

- ▶ Foundation of the least squares, maximum likelihood, ...

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A more sophisticated example

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$$\Theta = \{\boldsymbol{\theta} \in \mathbb{R}^d : \|\boldsymbol{\theta}\|_0 \leq s_0\}.$$

Theorem (Donoho, Johnstone 1990s)

If $s_0/d \rightarrow 0$, then

$$R_n^*(\Theta) = \frac{2s_0}{n} \log(d/s_0) \cdot (1 + o(1)).$$

- ▶ Key role in compressed sensing, sparse learning, ...

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What is this talk about?

$$R_n^{\text{Poly}}(\Theta) = \inf_{\hat{\theta}(\cdot) \in \text{Poly}} \sup_{\theta_0 \in \Theta} \mathbb{E}_{\theta_0} L(\hat{\theta}(x_1^n), \theta_0)$$

Developments

- ▶ Often we expect $R_n^{\text{Poly}}(\Theta) \geq R_n(\Theta)$
- ▶ Accurate predictions
- ▶ Convergence of fields (Statistics, CS Theory, Physics)
- ▶ New algorithms?

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Outline

1 Local algorithms

2 Landscapes

3 SDP relaxations

4 Conclusion

Local and message passing algorithms

Example #1: Sparse low-rank matrix

Unknowns:

$$\Theta(s_0, d) = \left\{ \boldsymbol{\theta} \in \{0, 1\}^d : \|\boldsymbol{\theta}\|_0 = s_0 \right\}.$$

Loss:

$$L(\hat{\boldsymbol{\theta}}, \boldsymbol{\theta}) = \frac{1}{d} \text{Hamming}(\hat{\boldsymbol{\theta}}, \boldsymbol{\theta})$$

Example #1: Sparse low-rank matrix

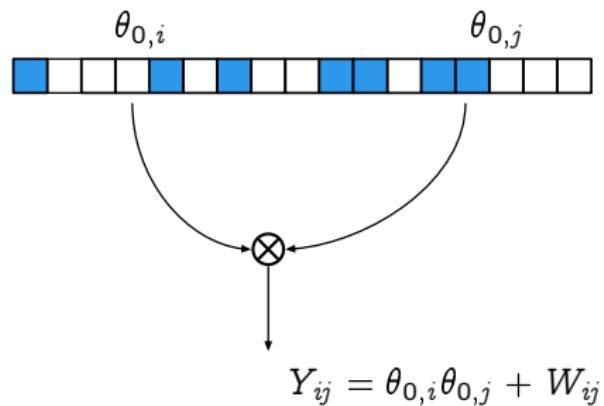
Data: (x_1, \dots, x_n)

$$x_\ell = (i_\ell, j_\ell, Y_{i_\ell, j_\ell}) \in [d] \times [d] \times \mathbb{R},$$

$$i_\ell, j_\ell \sim_{iid} \text{Unif}([d]),$$

$$Y_{i_\ell, j_\ell} |_{i_\ell, j_\ell} \sim N(\theta_{0, i_\ell} \theta_{0, j_\ell}, \sigma^2)$$

Pictorially



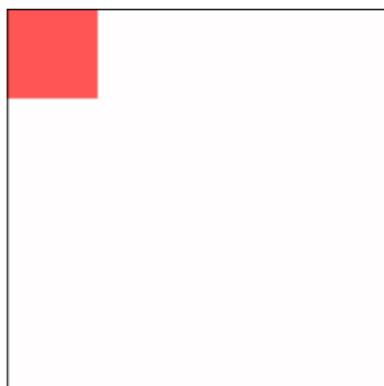
Example #1: A different formulation

Data $\mathbf{Y} \in \mathbb{R}^{n \times n}$

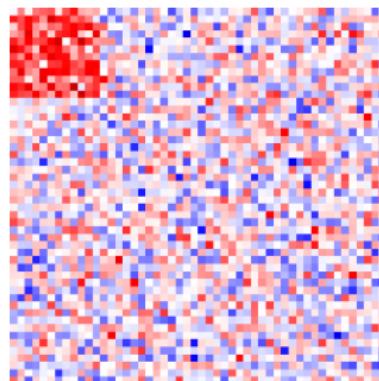
$$\mathbf{Y} = \mathcal{P}_E(\boldsymbol{\theta}_0 \boldsymbol{\theta}_0^\top + \mathbf{W})$$

- ▶ $E \subseteq \binom{[d]}{2}$ uniformly random s.t. $|E| = n$
- ▶ $\mathcal{P}_E : \mathbb{R}^{d \times d} \rightarrow \mathbb{R}^{d \times d}$: projector that zeroes entries not in E
- ▶ $(W_{ij})_{i < j} \sim_{iid} N(0, \sigma^2)$, $\mathbf{W} = \mathbf{W}^\top$

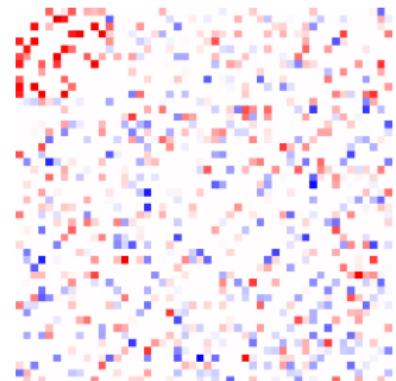
$$\theta_0 \theta_0^T$$



$$\theta_0 \theta_0^T + W$$



$$Y = \mathcal{P}_E(\theta_0 \theta_0^T + W)$$



Example #1: Yet another formulation

- ▶ $G = (V, E) \sim \mathcal{G}(d, n)$
(uniform random with d vertices, n edges)
- ▶ For each $(i, j) \in E$

$$Y_{ij} = \theta_i \theta_j + W_{ij}, \quad W_{ij} \sim N(0, \sigma^2)$$

Problem parameters

$$\delta = \frac{n}{d} \quad (\text{half}) \text{ average graph degree}$$

$$\varepsilon = \frac{s_0}{d} \quad \text{sparsity}$$

Asymptotics

- ▶ Dense graph: $n \asymp d^2$, $s_0 \asymp \sqrt{d}$ (\sim Planted clique problem)
- ▶ Sparse graph: $n \asymp d$, $s_0 \asymp d$ (δ, ε fixed)

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$$R^{\text{Poly}}(\delta, \varepsilon; d) \equiv R_{n=d\delta}^{\text{Poly}}(\Theta(s_0 = d\varepsilon, d))$$

$$\lim_{d \rightarrow \infty} R^{\text{Poly}}(\delta, \varepsilon; d) = ?$$

We do not know, but ...

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First simplifications

- ▶ Worst case prior

$$\theta \sim \text{Unif}(\Theta(s_0, d))$$

- ▶ Roughly ($\varepsilon = s_0/d$)

$$(\theta_i)_{i \leq d} \sim_{iid} \text{Bern}(\varepsilon)$$

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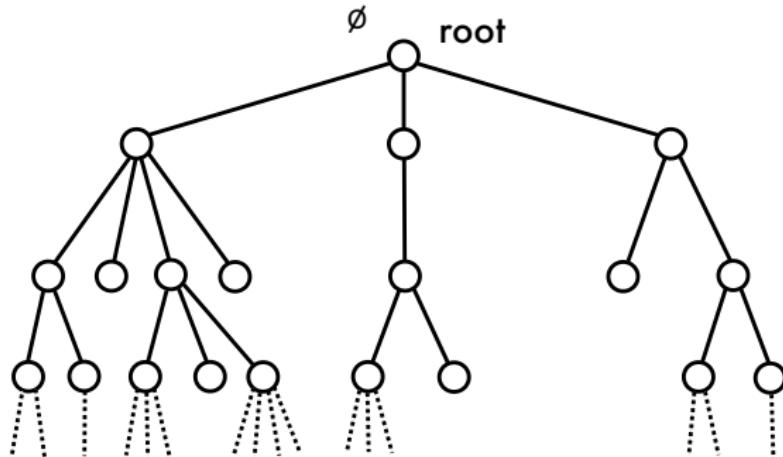
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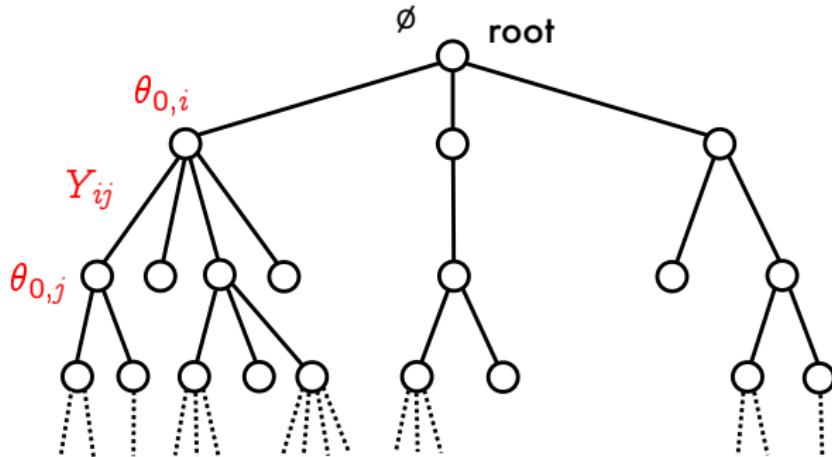
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Local weak limit: $d \rightarrow \infty$



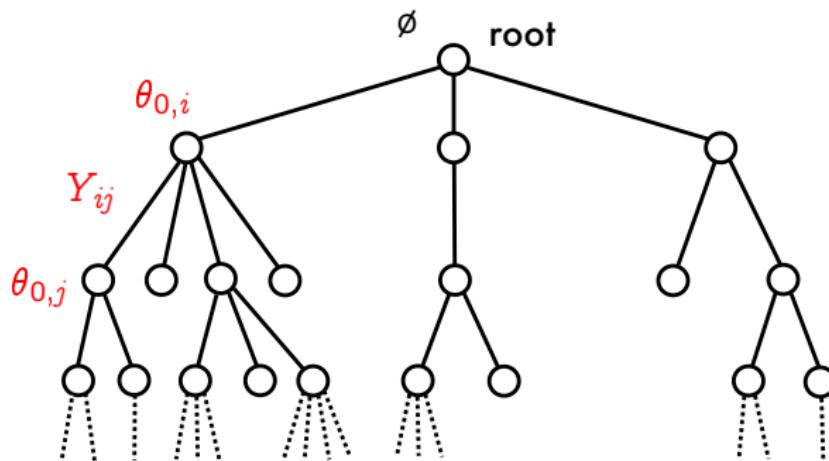
$$G \xrightarrow{lwc} \text{GW}(\text{Pois}(2\delta))$$

Local weak limit: $d \rightarrow \infty$



$$(G, \theta, Y) \xrightarrow{lwc} (\text{GW}(\text{Pois}(2\delta)), \theta, Y)$$

First guess



$$\lim_{d \rightarrow \infty} R^{\text{Poly}}(\delta, \varepsilon, d) \stackrel{?}{=} \inf_{\hat{\theta}} \mathbb{P}_{\text{Tree}}(\hat{\theta}_\emptyset(Y) \neq \theta_{0,\emptyset})$$

It gets interesting

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How do you define the r.h.s.?

A natural idea:

$$Y_\ell^0 = (Y_{ij} : d(\emptyset, i) \leq \ell),$$
$$\hat{\theta}_\emptyset(Y_\ell^0) = \arg \max_{\sigma \in \{0,1\}} \mathbb{P}(\theta_\emptyset = \sigma | Y_\ell^0).$$

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Risk of local algorithms

$$R^{\text{loc}}(\delta, \varepsilon) = \lim_{\ell \rightarrow \infty} \mathbb{P}_{\text{Tree}}(\hat{\theta}_\emptyset(Y_\ell^0) \neq \theta_{0,\emptyset})$$

Remark

Belief propagation achieves $R^{\text{loc}}(\delta, \varepsilon)$

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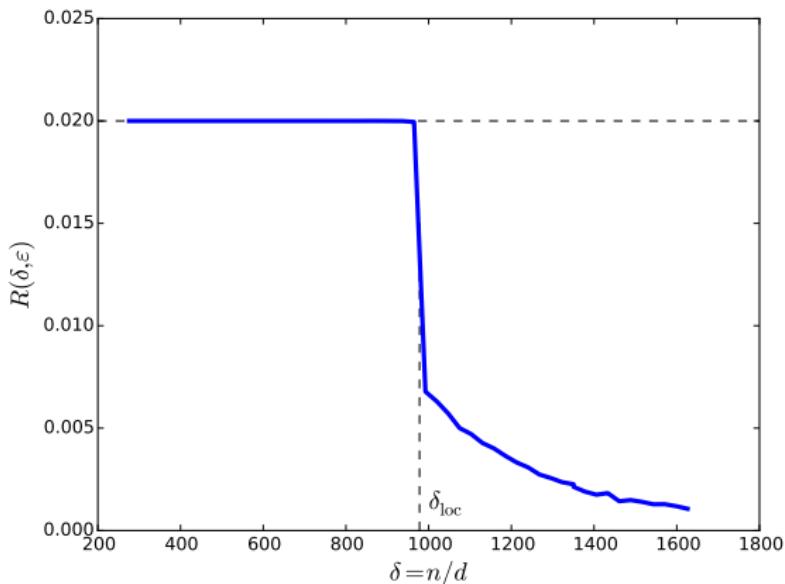
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Belief propagation achieves $R^{\text{loc}}(\delta, \varepsilon)$

Can be computed

($\varepsilon = 0.02, \sigma = 1.5$)



$$R^{\text{loc}}(\delta, \varepsilon) = \lim_{\ell \rightarrow \infty} \mathbb{P}_{\text{Tree}}(\hat{\theta}_\emptyset(Y_\ell^0) \neq \theta_{0,\emptyset})$$

Phase transition for local algorithms

Theorem (\sim Deshpande, Montanari, 2015)

As $\varepsilon \rightarrow 0$ (with $\delta \rightarrow \infty$, σ fixed)

$$\frac{1}{\varepsilon} R^{\text{loc}}(\delta, \varepsilon) \rightarrow \begin{cases} 1 & \text{if } \delta \leq (1 - o_\varepsilon(1)) \cdot \frac{\sigma^2}{2e\varepsilon^2}, \\ 0 & \text{if } \delta \geq (1 + o_\varepsilon(1)) \cdot \frac{\sigma^2}{2e\varepsilon^2}, \end{cases}$$

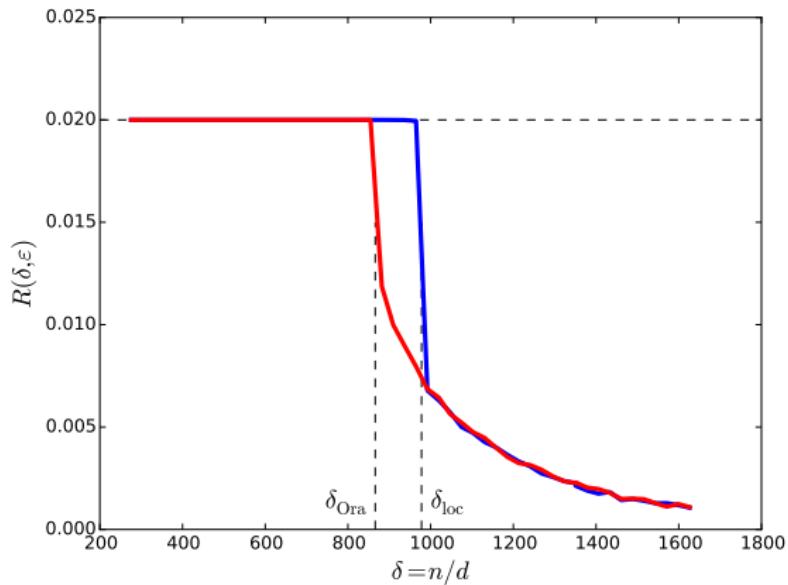
A different definition of the limit

$$Y_\ell^+ = \{ (Y_{ij} : d(\emptyset, i) \leq \ell); (\theta_i : d(\emptyset, i > \ell)) \}$$

$$\hat{\theta}_\emptyset(Y_\ell^+) = \arg \max_{\sigma \in \{0,1\}} \mathbb{P}(\theta_\emptyset = \sigma | Y_\ell^+)$$

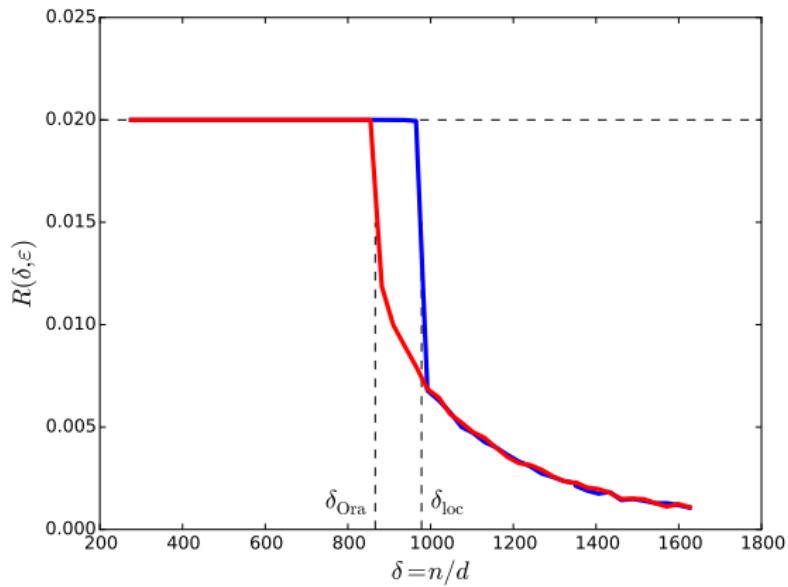
$$R^{\text{Ora}}(\delta, \varepsilon) = \lim_{\ell \rightarrow \infty} \mathbb{P}_{\text{Tree}}(\hat{\theta}_\emptyset(Y_\ell^0) \neq \theta_{0,\emptyset})$$

It also can be computed



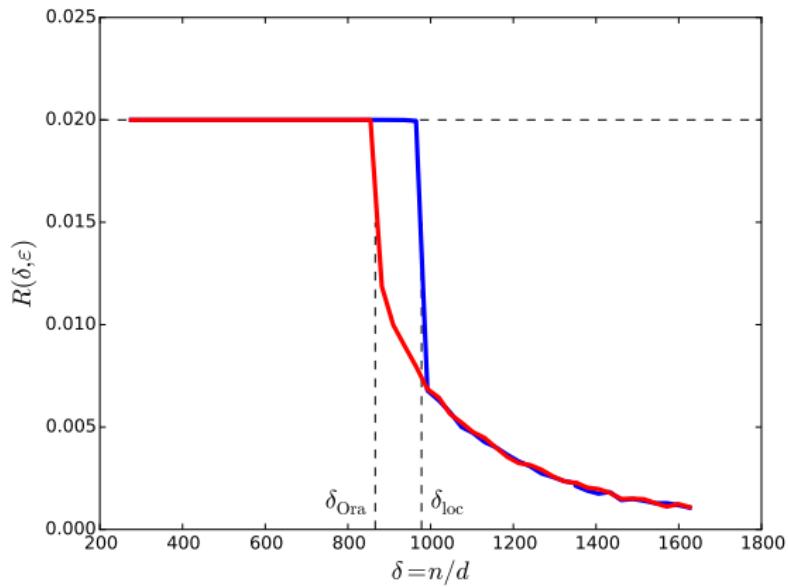
$$R^{\text{Ora}}(\delta, \varepsilon) \leq R^*(\delta, \varepsilon) \leq R^{\text{Poly}}(\delta, \varepsilon) \leq R^{\text{loc}}(\delta, \varepsilon)$$

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Minimax risk

Theorem (\sim Montanari, 2015)

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- ▶ Sharp threshold (dense): Lelarge, Miolane 2017; Barbier et al. 2017
- ▶ Do not know of any polytime algorithm working for

$$1.01 \cdot \frac{\sigma^2}{2\varepsilon} \log(1/\varepsilon) < \delta < 0.99 \cdot \frac{\sigma^2}{2e\varepsilon^2}.$$

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Open problems

- ▶ Can we beat $R^{\text{loc}}(\delta, \varepsilon)$ by Gibbs sampling?
- ▶ Can we beat $R^{\text{loc}}(\delta, \varepsilon)$ by convex optimization?
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Ubiquitous

Sparse principal component analysis

- ▶ Berthet, Rigollet, 2013
- ▶ Deshpande, Montanari, 2014
- ▶ Barbier et al 2016; Miolane 2017

Hidden clique problem

(2 asymmetric communities)

- ▶ Jerrum, 1992
- ▶ Feige, Krauthgamer, 2000
- ▶ Deshpande, Montanari, 2015; Montanari 2015

Community detection ($k \geq 5$ symmetric communities)

- ▶ Decelle, Krzakala, Moore, Zdeborova 2011
- ▶ Bordenave, Lelarge, Massouline 2015
- ▶ Abbe, Sandon 2015

Tensor PCA

- ▶ See below

Landscapes

Empirical risk minimization/M-estimation

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Rationale

$$\theta_0 = \arg \min_{\theta \in \mathbb{R}^d} \mathcal{L}(\theta), \quad \mathcal{L}(\theta) = \mathbb{E} \hat{\mathcal{L}}_n(\theta)$$

- ▶ What can we say generically?
- ▶ How does complexity show up?

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Uniform convergence

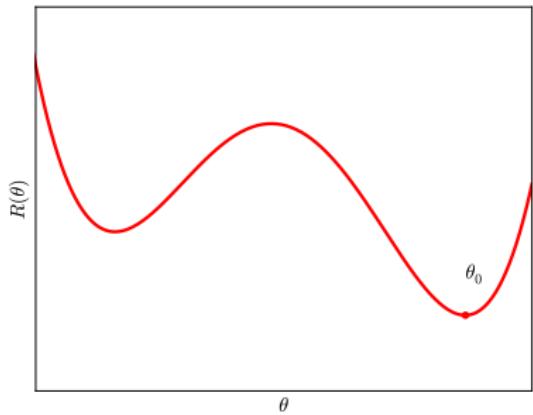
Theorem (Vapnik, Chervonenkis, 1968; ...)

Under conditions [omitted], with high probability

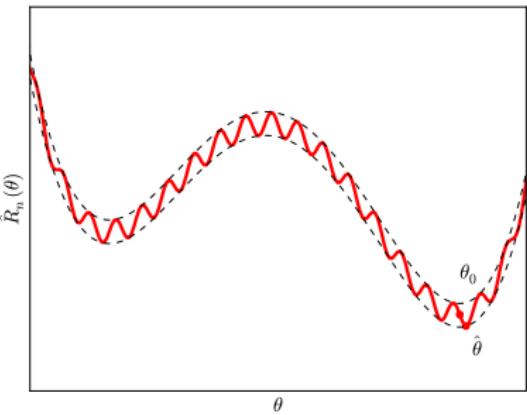
$$\sup_{\theta \in \Theta} |\widehat{\mathcal{L}}_n(\theta) - \mathcal{L}(\theta)| \leq C \sqrt{\frac{d_*}{n}}.$$

(d_ = VC dimension; ...).*

Uniform convergence



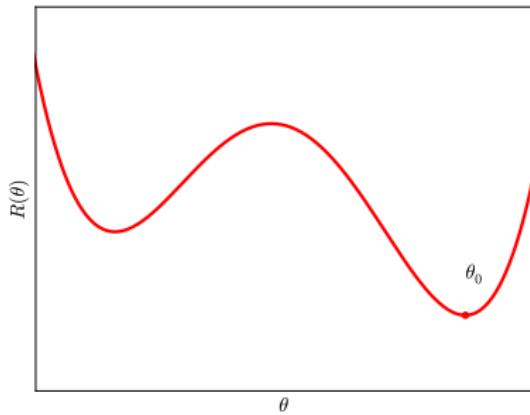
Population risk



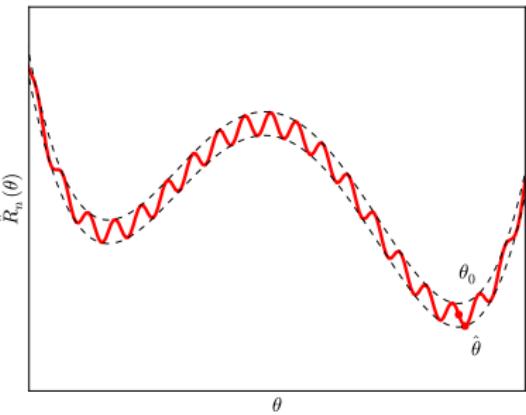
Empirical risk

Will optimization algorithms get stuck in local minima?
Landscape analysis

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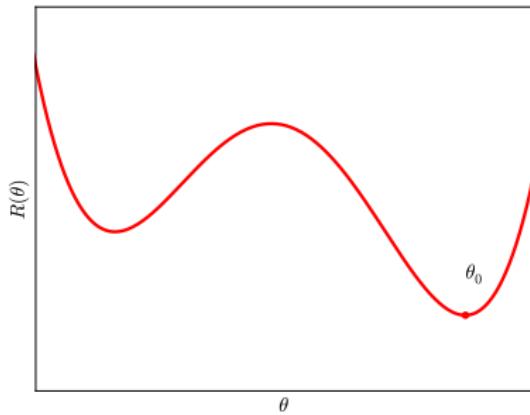
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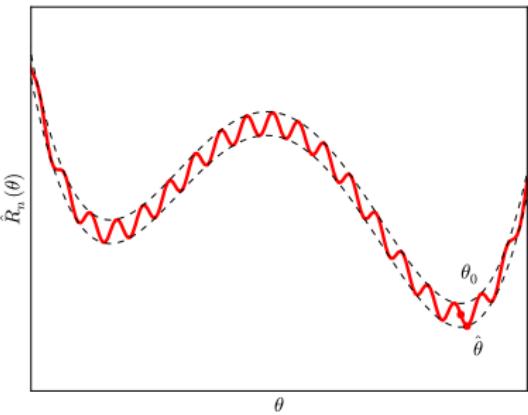
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Assumptions

$\theta \in B^p(r) =$ Ball of radius r in \mathbb{R}^p

Data: Z_1, Z_2, \dots, Z_n iid

A1 $\nabla_\theta \ell(\theta; Z)$ is τ^2 -sub-Gaussian

A2 For any $\lambda \in B^p(1)$, $\mathcal{Z}_\lambda \equiv \langle \lambda, \nabla^2 \ell(\theta; Z) \lambda \rangle$ is τ^2 -sub-Exponential.

A3 The Hessian of the population risk at $\mathbf{0}$ is bounded by a polynomial

$$\|\nabla^2 \mathcal{L}(\mathbf{0})\|_{\text{op}} \leq \tau^2 d^C.$$

A4 The Hessian of the loss is Lipschitz continuous with integrable constant

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'Under mild assumptions'

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A2 For any $\lambda \in B^p(1)$, $\mathcal{Z}_\lambda \equiv \langle \lambda, \nabla^2 \ell(\theta; Z) \lambda \rangle$ is τ^2 -sub-Exponential.

A3 The Hessian of the population risk at $\mathbf{0}$ is bounded by a polynomial

$$\|\nabla^2 \mathcal{L}(\mathbf{0})\|_{\text{op}} \leq \tau^2 d^C.$$

A4 The Hessian of the loss is Lipschitz continuous with integrable constant

$$\mathbb{E}\{\|\nabla^2 \ell(\cdot; Z)\|_{\text{Lip}}\} \leq \tau^3 d^C.$$

'Under mild assumptions'

Assumptions

$\theta \in B^p(r) =$ Ball of radius r in \mathbb{R}^p

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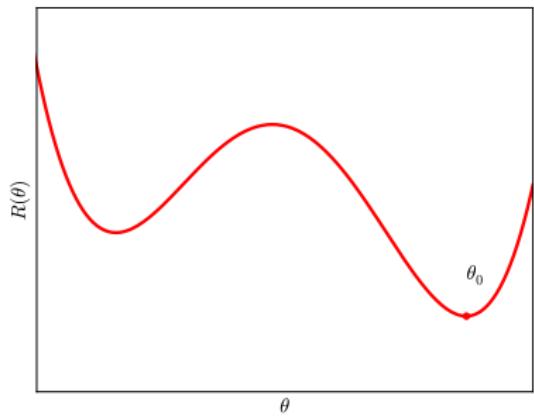
Lemma

Under assumptions A1, A2, A3, A4, if $n \geq Cp \log p$, then with probability at least $1 - \delta$, the following hold:

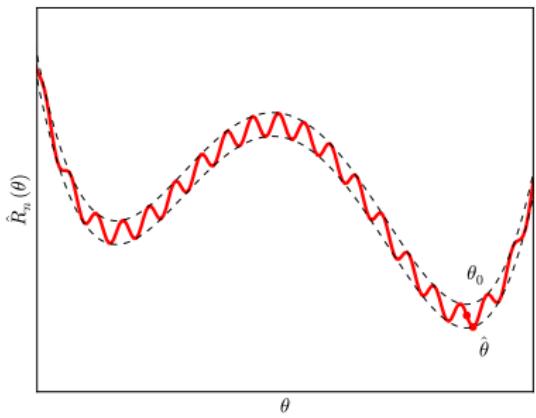
$$\sup_{\theta \in \mathbb{B}^p(r)} \left\| \nabla \widehat{\mathcal{L}}_n(\theta) - \nabla \mathcal{L}(\theta) \right\|_2 \leq \tau \sqrt{\frac{Cd \log n}{n}},$$

$$\sup_{\theta \in \mathbb{B}^p(r)} \left\| \nabla^2 \widehat{\mathcal{L}}_n(\theta) - \nabla^2 \mathcal{L}(\theta) \right\|_{\text{op}} \leq \tau^2 \sqrt{\frac{Cd \log n}{n}}.$$

This cannot happen!

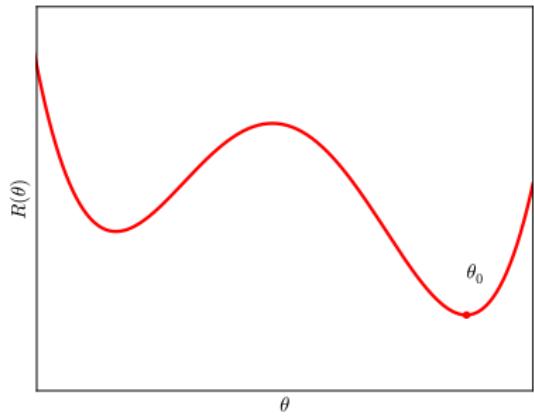


Population risk

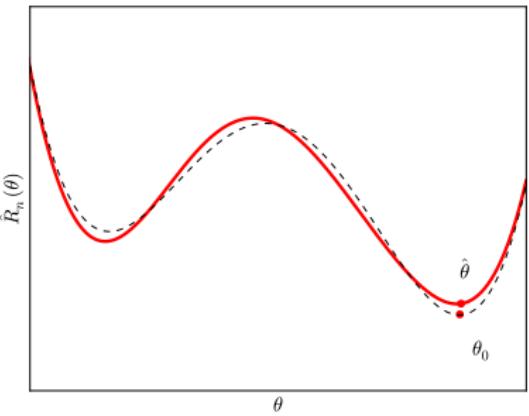


Empirical risk

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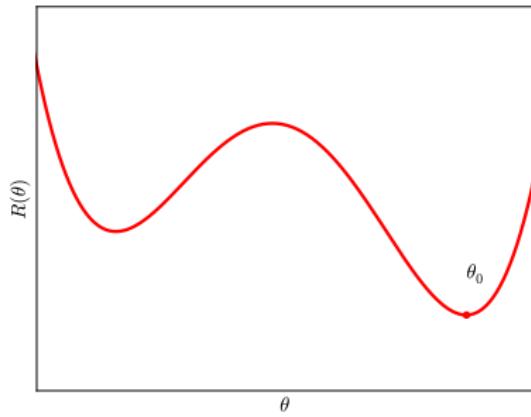
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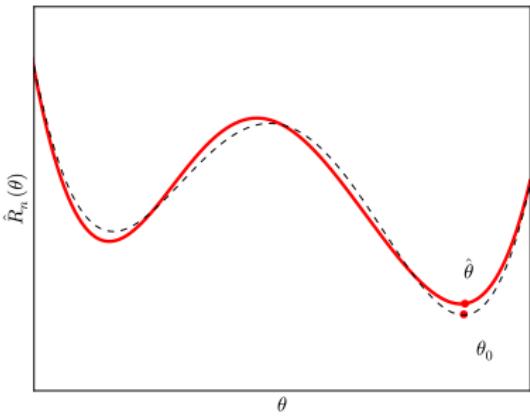
Empirical risk

Nice population risk \Rightarrow Nice empirical risk

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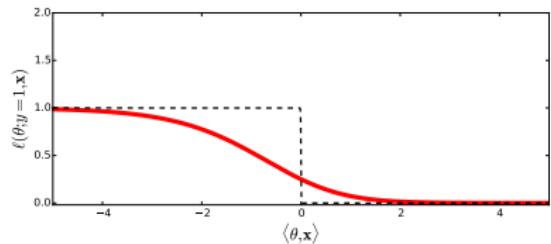
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Example: Binary classification



$z_i = (y_i, \mathbf{x}_i)$, $y_i \in \{0, 1\}$, $\mathbf{x}_i \in \mathbb{R}^d$

$$\mathbb{P}(y_i = 1 | \mathbf{x}_i) = \sigma(\langle \boldsymbol{\theta}_0, \mathbf{x}_i \rangle)$$

$$\widehat{\mathcal{L}}_n(\boldsymbol{\theta}) = \frac{1}{n} \sum_{i=1}^n (y_i - \sigma(\langle \boldsymbol{\theta}, \mathbf{x}_i \rangle))^2.$$

- ▶ Rosenblatt 1958 (Perceptron); ... many extensions
- ▶ More robust than logistic regression

Sample application: Binary classification

Theorem (Mei, Bai, Montanari 2017)

Assume X_i to be centered, sub-Gaussian, with $\mathbb{E}\{XX^\top\} \succeq \delta I_d$.
For nice^a functions σ , whp:

1. The population risk has a unique critical point $\hat{\theta}_n$.
2. Gradient descent converges exponentially fast to $\hat{\theta}_n$.
3. The estimation error is $\|\hat{\theta}_n - \theta_0\|_2 \leq C \sqrt{(d \log n)/n}$.

^a $\sigma'(x) > 0$, $\|\sigma'\|_\infty, \|\sigma''\|_\infty, \|\sigma'''\|_\infty \leq C$.

Similar results for robust regression, one-bit compressed sensing, ...
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Non-convex literature

Convergence to ‘statistical neighborhood’

- ▶ Loh, Wainwright, 2012
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- ▶ Yang, Wang, Liu, Eldar, Zhang, 2015
- ▶ ...

Smart initialization

- ▶ Keshavan, Montanari, Oh, 2009
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Unique local minimum

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- ▶ What can we say generically?
- ▶ How does complexity show up?

Intuition

Population risk very flat \Leftrightarrow Many local minima

[Close to where local algorithms fail?]

Simplest example: Spiked tensor model

- ▶ Unknown parameter $\theta_0 \in \mathbb{R}^n$, $\|\theta_0\|_2 = 1$
- ▶ Data¹

$$Y = \lambda \theta_0^{\otimes k} + W$$

- ▶ W = symmetric Gaussian noise tensor.
 $(W_{i_1, \dots, i_k})_{i_1 < \dots < i_k} \sim_{iid} N(0, 1/n)$

¹Equivalently, $Y_{i_1, \dots, i_k} = \lambda \theta_{0,i_1} \cdots \theta_{0,i_k} + W_{i_1, \dots, i_k}$.

Spiked tensor model: What do we know?

$$Y = \lambda \theta_0^{\otimes k} + W$$

Theorem (Montanari, Richard, 2014; Hopkins, Shi, Steurer, 2015)

For any $\varepsilon > 0$, there exist constants λ_{IT} , $\lambda_{\text{ML}}(\varepsilon)$, $C(\varepsilon)$, such that:

- ▶ If $\lambda > \lambda_{\text{ML}}(\varepsilon)$, then $\mathbb{E}\{|\langle \hat{\theta}^{\text{ML}}, \theta_0 \rangle|\} \geq 1 - \varepsilon$.
- ▶ No estimator can achieve $\mathbb{E}\{|\langle \hat{\theta}, \theta_0 \rangle|\} \geq \varepsilon$ unless $\lambda > \lambda_{\text{IT}}$.
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No efficient estimator is known for $1 \ll \lambda \ll n^{(k-2)/4}$!

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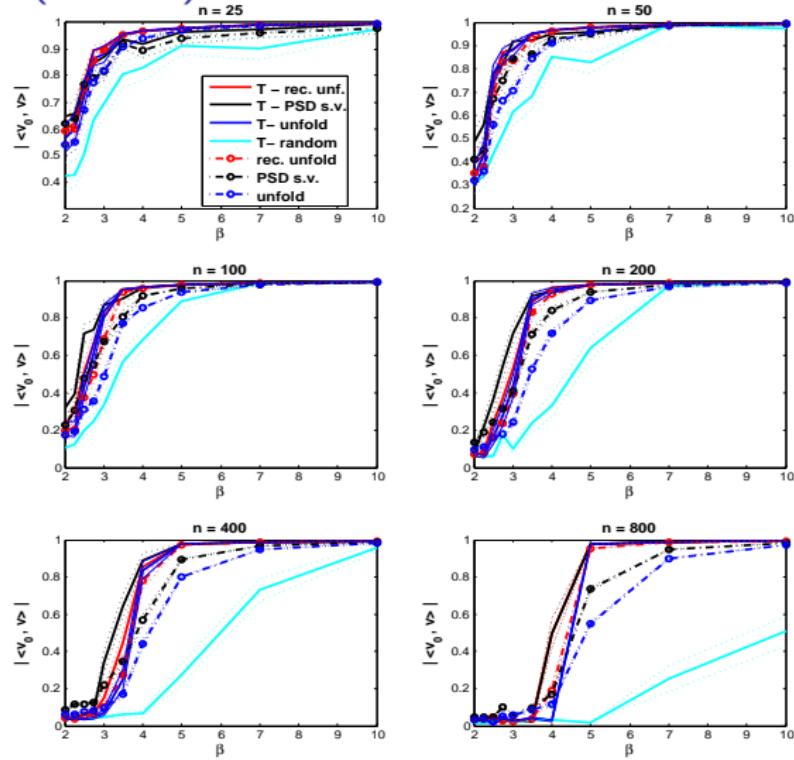
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More precise results

- ▶ Montanari, Reichman, Zeitouni, 2015
- ▶ Bandeira, Perry, Wein, 2017
- ▶ Krzakala, Lelarge, Miolane, Zdeborova, 2017
- ▶ ...

In practice ($k = 3$)



What does the landscape look like?

Maximum likelihood

$$\begin{aligned} & \text{minimize} && \hat{\mathcal{L}}_n(\boldsymbol{\theta}) = \| \mathbf{Y} - \lambda \boldsymbol{\theta}^{\otimes k} \|_F \\ & \text{subject to} && \|\boldsymbol{\theta}\|_2 = 1 \end{aligned}$$

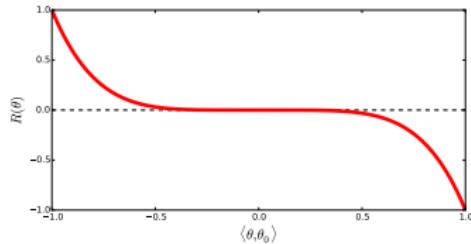
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Risk



Maximum likelihood

$$\begin{aligned} \text{minimize } & \widehat{\mathcal{L}}_n(\boldsymbol{\theta}) = -\langle \mathbf{Y}, \boldsymbol{\theta}^{\otimes k} \rangle \\ \text{subject to } & \|\boldsymbol{\theta}\|_2 = 1 \end{aligned}$$

‘Population’ risk

$$\mathcal{L}(\boldsymbol{\theta}) = -\lambda \langle \boldsymbol{\theta}_0, \boldsymbol{\theta} \rangle^k$$

Back-of-the-envelope

Expected gradient

$$\nabla \mathcal{L}(\boldsymbol{\theta}) = -k\lambda \langle \boldsymbol{\theta}, \boldsymbol{\theta}_0 \rangle^{k-1} \boldsymbol{\theta}_0$$

Random initialization $\langle \boldsymbol{\theta}, \boldsymbol{\theta}_0 \rangle = \Theta(n^{-1/2})$:

$$\begin{aligned}\langle \boldsymbol{\theta}_0, \nabla \widehat{\mathcal{L}}_n(\boldsymbol{\theta}) \rangle &= -k\lambda \langle \boldsymbol{\theta}, \boldsymbol{\theta}_0 \rangle^{k-1} - k \langle \mathbf{W}, \boldsymbol{\theta}_0 \otimes \boldsymbol{\theta}^{\otimes(k-1)} \rangle \\ &= -\lambda \Theta(n^{-(k-1)/2}) + \Theta(n^{-1/2})\end{aligned}$$

- ▶ Convergence: $\lambda \gg n^{(k-2)/2}$

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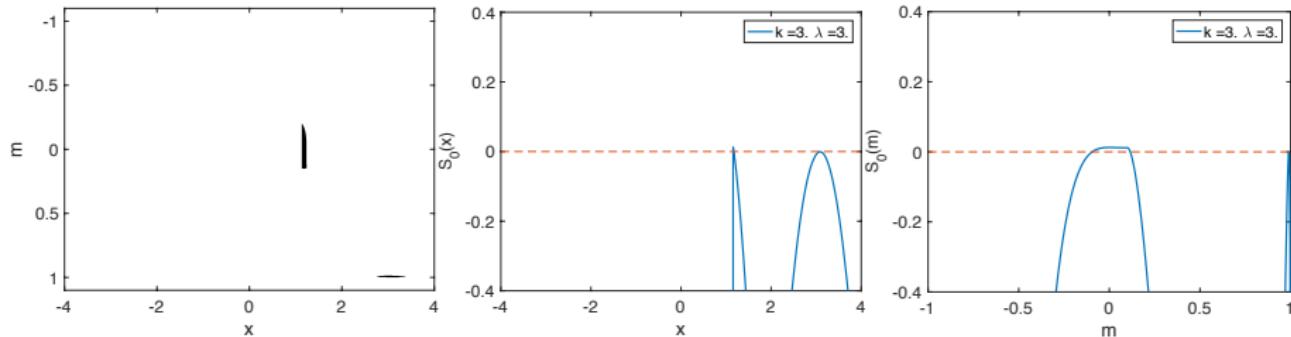
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- Convergence: $\lambda \gg n^{(k-2)/2}$

Expected number of local minima: $k = 3, \lambda = 3$



- Exponential in black region $(m = \langle \theta, \theta_0 \rangle, x = \langle Y, \theta^{\otimes k} \rangle)$
- N = number of local minima

$$\mathbb{E}N(m, x) = e^{nS_0(m, x) + o(n)}$$

[Ben Arous, Mei, Montanari, Nica, 2017]

Complexity of landscape $\overset{?}{\leftrightarrow}$ Complexity for local algorithms

SDP relaxations

An emerging dichotomy

In several statistical estimation problems

- ▶ Either local (or message-passing) algorithms work . . .
- ▶ . . . or SDP hierarchies do not work

Why?

An emerging dichotomy

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Why?

A possible explanation

Perhaps SDPs on random instances can be solved by local algorithms...

Simplest example

Centered adjacency matrix of $G = (V, E)$ (d = average degree)

$$A_{ij}^{\text{cen}} = \begin{cases} 1 - \frac{d}{n} & \text{if } (i, j) \in E, \\ -\frac{d}{n} & \text{otherwise.} \end{cases}$$

SDP(A^{cen}):

$$\begin{aligned} & \text{maximize} && \langle A^{\text{cen}}, X \rangle, \\ & \text{subject to} && X \in \mathbb{R}^{n \times n}, \quad X \succeq 0, \\ & && X_{ii} = 1. \end{aligned}$$

- ▶ Graph clustering, embedding, testing latent structure,...

What does it mean?

Input: Graph $G_n = (V_n, E_n)$

1. Generate $z = (z(i))_{i \in V} \sim_{iid} N(0, 1)$
2. Compute, for each $v \in V_n$, $\xi_v = F(B_\ell(v; G_n), z|_{B_\ell(v; G_n)})$
3. Output $X = \mathbb{E}_z\{\xi\xi^\top\} + \dots \in \mathbb{R}^{n \times n}$

Can this achieve $\langle A^{\text{cen}}, X \rangle \geq (1 - o_n(1))\text{SDP}(A^{\text{cen}})$?

Erdős-Renyi random graphs

Theorem (Fan, Montanari, 2016)

Let $G \sim \mathcal{G}(n, d/n)$ and $\mathbf{A}^{\text{cen}} = \mathbf{A}_G^{\text{cen}}$. Then, a.s.,

$$\begin{aligned} 2\sqrt{d} \left(1 - \frac{1}{d+1}\right) &\leq \liminf_{n \rightarrow \infty} \frac{1}{n} \text{SDP}(\mathbf{A}^{\text{cen}}) \leq \\ &\leq \limsup_{n \rightarrow \infty} \frac{1}{n} \text{SDP}(\mathbf{A}^{\text{cen}}) \leq 2\sqrt{d} \left(1 - \frac{1}{2d}\right). \end{aligned}$$

Further, the lower bound is achieved by local algorithms.

- ▶ A local algorithm achieves $8/9$ of $\text{SDP}(\mathbf{A}^{\text{cen}})$.

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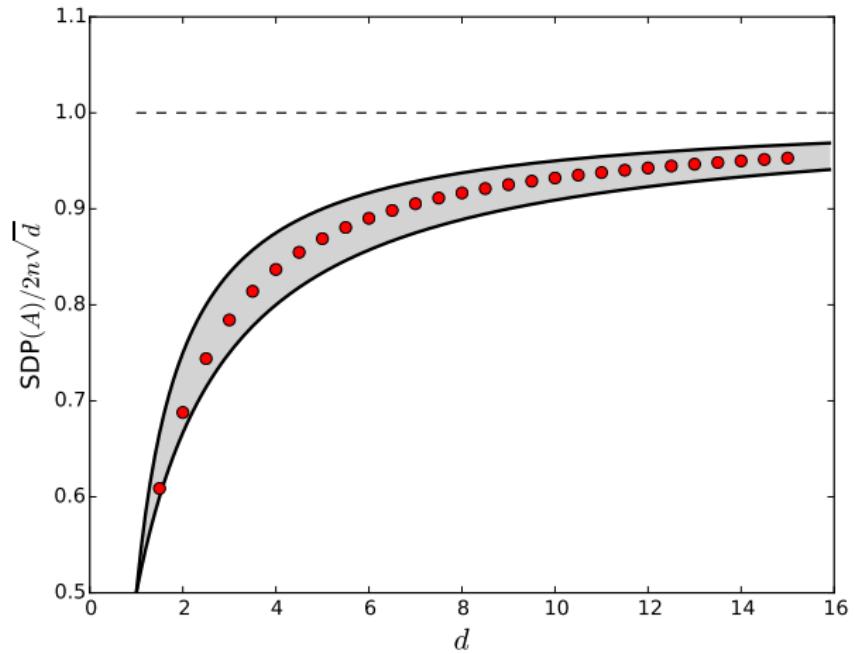
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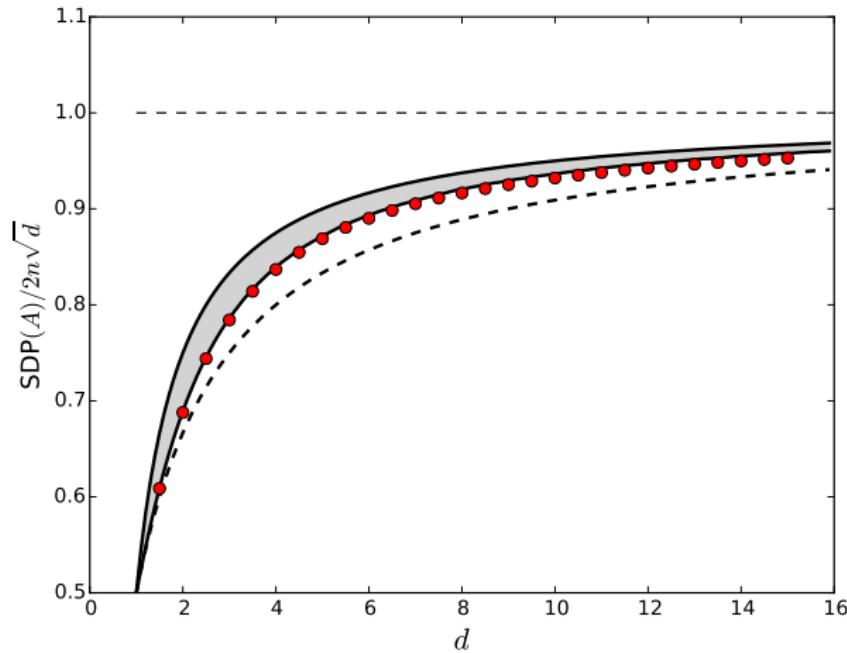
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Bounds vs numerical simulations



[Related results for planted partition model, regular graphs, ...]

A better local algorithm (see paper)



[Related results for planted partition model, regular graphs, ...]

Conclusion

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- ▶ Which statistical problems are tractable?
- ▶ Multiple points of view:
 - ▶ Local-message passing algorithms
 - ▶ Landscape analysis
 - ▶ SDP hierarchies

Thanks!

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