You have 90 minutes for this examination; the number of points for each question indicates roughly how many minutes you should spend on that question. Make sure you print your name and sign the Honor Code statement below. During this exam you are allowed to consult two double-sided pages of notes that you have prepared ahead of time; other than that, you may not consult books, notes, your laptop or cell phone, or any other materials. If there is a trivial detail that you need for one of your answers but cannot recall, you may ask the course staff for help.

Note: be sure to write only on the front sides of pages. We’ll be using Gradescope to grade the exams and may not see information written on the back sides. There are extra pages at the end of the exam, if you need more space for an answer.

I acknowledge and accept the Stanford University Honor Code. I have neither given nor received aid in answering the questions on this examination, and I have not consulted any external information other than two double-sided pages of prepared notes.

_______________________________________
(Signature)
_______________________________________
(Print your name, legibly!)
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Problem 1 (8 points)
Indicate whether each of the following statements is true or false, and explain your answer briefly.

(a) Multiple threads can execute in the same critical section simultaneously, as long as they run on different cores.

(b) MacOS just updated itself last night, and the only change in the update is a new processor scheduling algorithm. As soon as the update installed, one of the applications that you use regularly—Adobe Reader—starts crashing. You should file a bug report with Adobe, not Apple.

(c) The 4.4 BSD scheduler (the one that measures CPU usage for each thread) is non-preemptive.

(d) In Project 3, the object file party_test.o contains code that invokes the meet method of the Party class; the code for the Party class is in object file party.o. Because of this, there will be a symbol table entry for the meet method in party_test.o.
Problem 2 (4 points)

(Note: this project is no longer part of the class)

In Project 1, when the shell starts a new command, it creates a new process for that command, with a single thread. Suppose that instead of creating a new process with fork and execvp, it created a new thread within its own process to run the command: give two reasons why this would be a bad idea.

Problem 3 (10 points)

The very first priority-based scheduler we discussed in lecture makes use of several round-robin queues—on the order of 5 or 6—to manage all threads eligible for processing time and schedule them to execute in an order that leads to the shortest average completion time, or something very close to it.

a) (2 points) What heuristic does the OS use to identify threads as CPU-bound?

b) (4 points) An I/O-bound thread, when eligible for CPU time, typically occupies a spot in the highest priority queue. What types of I/O-bound threads allow for CPU-bound ones executing with lower priority to get processor time? In what situations would a CPU-bound thread be more likely to starve?

c) (4 points) Multi-queue schedulers typically set the time slice duration for lower priorities to be higher. For example, the time slice duration granted to threads with the highest priority might be, say, 4ms, whereas the time slice duration granted to threads running with the second highest priority might be 8ms (third-highest priority threads? 16ms, etc.) Give two reasons why higher-priority threads should be granted smaller time slices whereas lower-priority ones should be granted larger time slices.
Problem 4 (12 points)

The figure below shows the address translation mechanism used for 64-bit mode in Intel x86 processors, which was discussed in lecture. It uses 4 levels of page map to translate 48 bits of virtual address to 52-bit physical addresses.

(a) (2 points) How many bytes are in each page map entry in this scheme (not all bits of a page map entry are necessarily used)?

(b) (10 points) A friend comes to you and claims that the number of levels of page map could be decreased by increasing the page size. Specifically, they claim that if the page size is increased by a factor of 4x, only 3 levels of page map will be needed to translate 48-bit virtual addresses. Is your friend right? Redraw the mechanism above with 4x larger pages and only 3 levels of page map; how many total bits of virtual address does the new mechanism translate/use?
Problem 5 (40 points)

Because of construction, a multi-lane highway must merge to a single lane. Caltrans has hired you to automate the merge point. To do this, you must define a C++ class `Merge` with a constructor and two methods:

```cpp
class Merge {
public:
    Merge(int num_lanes);
    void wait(int lane);
    void merged();
private:
    // You must fill in the part below.
};
```

The `num_lanes` argument to the constructor indicates the total number of lanes that will be merging to a single lane. When a car arrives at the merge point it will invoke the `wait` method; the `lane` argument indicates the lane number in which the car arrives (this value will always be between 0 and `num_lanes-1`, inclusive). The `wait` method must not return until it is this car’s turn to pass the merge point into the single lane. Once the car has finished merging into the one lane, it will invoke the `merged` method; this indicates that this car has finished merging and it is safe for the next car to begin merging. The car threads coordinate among themselves using these two methods: there is not a separate flag-person thread.

Your solution must meet the following requirements:

- There is a single outgoing lane, and only one car can merge at a time (i.e., only one car can have returned from `wait` but not have called `merged` yet).
- If there are many cars waiting, they must take turns in order of lane (one car from each lane with a waiting car must merge before a second car merges from any lane).
- There will not necessarily be cars waiting in all (or any) lanes at any given point in time.
- Unlike a real highway, **you do not need to guarantee perfect FIFO order** for the cars in a single lane.
- You must write your solution in C++, using locks and condition variables.
- Use the monitor style for your code, as discussed in class and required for Project 3.
- Your solution must not use busy-waiting and must not allow starvation (you may assume that locks and condition variables will not starve threads).
- Try to make your solution simple and obvious; we will reduce your score by up to 20% if your solution is overly complicated.
Write the methods of the Merge class here:
Additional working space for any problem, if needed
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