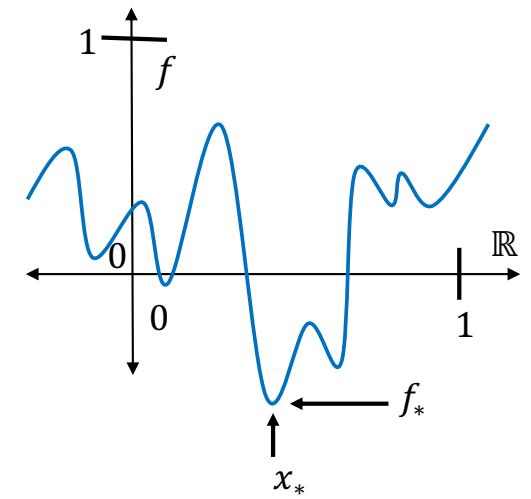
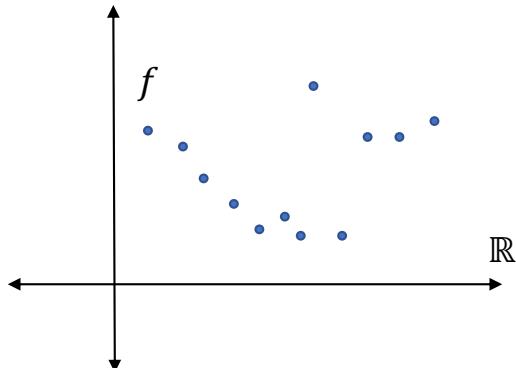


# Introduction to Optimization Theory

Lecture #16 - 11/10/20

MS&E 213 / CS 2690

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# Plan for Today

Recap

- Online linear optimization methods
- Follow the regularized leader (FTRL)
- Mirror decent

Mirror  
Descent

- Analysis
- Perspectives
- Comparison

SGD / SVRG

- Finite sum optimization
- Stochastic variance reduced gradient
- Minimizing sum of smooth functions

Dimension  
Dependent

Cutting plane and interior point methods

# Online Linear Optimization

- Given closed convex  $S \subseteq \mathbb{R}^n$
- Iterate for round  $t \in [T]$ 
  - We pick  $x_t \in S$
  - Adversary reveals penalty  $p_t \in \mathbb{R}^n$

## Goal:

- Minimize  $\text{regret}(T) \stackrel{\text{def}}{=} \sum_{t \in [T]} p_t^\top x_t - \min_{x \in S} \sum_{t \in [T]} p_t^\top x$
- Want average regret  $\rightarrow 0 : \lim_{T \rightarrow \infty} \frac{\text{regret}(T)}{T} \rightarrow 0$
- More generally: bound  $\sum_{t \in [T]} p_t^\top (x_t - z)$  for all  $z$

# Follow the Regularized Leader (FTRL)

## Online Linear Optimization

- For  $t \in [T]$
- We pick  $x_t \in S$
- Adversary picks  $p_t \in \mathbb{R}^n$
- $\sum_{t \in [T]} p_t^\top (x_t - z)$ .

## Algorithm

- “regularizer”  $r: S \rightarrow \mathbb{R}$  that is differentiable and  $\mu$ -strongly convex with respect to  $\|\cdot\|$
- Let  $\Phi_T(x) \stackrel{\text{def}}{=} \eta \sum_{t \in [T]} p_t^\top x + r(x)$
- Let  $\Phi_0(x) \stackrel{\text{def}}{=} r(x)$
- Let  $x_{T+1} = \operatorname{argmin}_{x \in S} \Phi_T(x)$

## Analysis

- $p_t^\top x_t \leq \frac{1}{\eta} [\Phi_t(x_{t+1}) - \Phi_{t-1}(x_t)] + \frac{\eta}{2\mu} \|p_t\|_*^2$
- $\sum_{t \in [T]} p_t^\top (x_t - z) \leq \frac{\eta}{2\mu T} \sum_{t \in [T]} \|p_t\|_*^2 + \frac{1}{\eta T} \left[ r(z) - \min_{x \in S} r(x) \right]$

# Another Online Learning Algorithm

## Idea

- Start at  $x_0$
- For  $t \in [T]$
- Take small step in “direction”  $-p_t$

## How Design Step?

- Pick strongly convex  $r$
- Center regularizer around  $x_t$ 
  - Notation:  $D(x||x_t)$
- Minimizer  $\eta \cdot p_t^T x + D(x||x_t)$

Mirror  
Descent

## Intuition

- *Similar to before*
- *$x_t$  do not move to much relative to reference point*

# Centered Regularizer?

## Bregman Divergence Distance

- Let  $r$  be differentiable and  $\mu$ -strongly with respect to some norm.
- $D_r(x||c) = r(x) - [r(c) + \nabla r(c)^\top(x - c)]$

## Notes

- $D_r(x||c) \geq \frac{\mu}{2} \|x - c\|^2$
- $D_r(x||c) = 0 \Leftrightarrow x = c$
- $D_r(x||y)$  not necessarily  $D_r(y||x)$

## Examples

- $r(x) = \frac{1}{2} \|x - d\|_2^2$ 
  - $D_r(x||c) = \frac{1}{2} \|x - c\|_2^2$
  - Doesn't depend on  $d$
- $r(x) = \sum_{i \in [n]} x_i \log x_i$ 
  - $D_r(x||c) = \sum_{i \in [n]} x_i \log x_i$
  - KL-divergence

# Bregman Projection

## Bregman Divergence Distance

- Let  $r$  be differentiable and  $\mu$ -strongly convex with respect to some norm.
- $D_r(x||c) = r(x) - [r(c) + \nabla r(c)^\top(x - c)]$

## Bregman Projection

- For closed convex  $S \subseteq \mathbb{R}^n$  let  $\pi_S^r(y) = \underset{x \in S}{\operatorname{argmin}} D_r(x||y)$

## Bregman Pythagorean Theorem (obtuse angles)

- $D_r(x||y) + D_r(y||z) \leq D_r(x||z)$  for all  $x \in S$
- $\Leftrightarrow y = \pi_S^r(z)$

## Proof

- Algebra and optimality of projection

# Mirror Descent Analysis

- $D_r(x||c) = r(x) - [r(c) + \nabla r(c)^\top(x - y)]$
- $x_{t+1} = \operatorname{argmin}_{x \in S} \eta \cdot p_t^\top x + D_r(x||x_t)$

## Lemmas

- $p_t^\top(x_t - z) \leq \frac{1}{\eta}[D_r(z||x_t) - D_r(z||x_{t+1})] + \frac{\eta}{2\mu} \|p_t\|_*^2$  (*Mirror Descent*)
- $p_t^\top x_t \leq \frac{1}{\eta}[\Phi_t(x_{t+1}) - \Phi_{t-1}(x_t)] + \frac{\eta}{2\mu} \|p_t\|_*^2$  (*FTRL*)

## Theorem: Mirror Descent

- $\sum_{t \in [T]} p_t^\top(x_t - z) \leq \sum_{t \in [T]} \frac{\eta}{2\mu} \|p_t\|_*^2 + \frac{1}{\eta} D_r(z||x_0)$

- *Similar bound!*
- *Resulting methods are often the same!*

## Theorem: Dual Averaging (FTRL)

- $\sum_{t \in [T]} p_t^\top(x_t - z) \leq \sum_{t \in [T]} \frac{\eta}{2\mu} \|p_t\|_*^2 + \frac{1}{\eta} \left[ r(z) - \min_{x \in S} r(x) \right]$

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# Mirror Descent Analysis

- $D_r(x||c) = r(x) - [r(c) + \nabla r(c)^\top(x - c)]$
- $x_{t+1} = \operatorname{argmin}_{x \in S} \eta \cdot p_t^\top x + D_r(x||x_t)$

## Lemma

- $p_t^\top(x_t - z) \leq \frac{1}{\eta} [D_r(z||x_t) - D_r(z||x_{t+1})] + \frac{\eta^2}{2\mu} \|p_t\|_*^2$  (*Mirror Descent*)

## Proof

yields “Pythagorean Theorem”

- $D_r(x||z) - D_r(x||y) - D_r(y||z) = \nabla D_r(y||z)^\top(x - y)$  for all  $x, y, z$
- $(\eta \cdot p_t + \nabla D_r(x_{t+1}||x_t))^\top(z - x_{t+1}) \geq 0$
- $\eta \cdot p_t^\top(x_t - z) \leq \eta \cdot p_t^\top(x_t - x_{t+1}) - D_r(x_{t+1}||x_t)$   
 $+ D_r(z||x_t) - D_r(z||x_{t+1})$
- $D_r(x_{t+1}||x_t) \geq \frac{\mu}{2} \|x_{t+1} - x_t\|^2$

# Another Perspective

- Gradient's are in "dual space"
- Want to map dual to primal to take step

## Mirror Map

- Let  $r: T \rightarrow \mathbb{R}$  be differentiable function with  $S \subseteq T$  such that for all  $g \in \mathbb{R}^n$  it is the case that  $\exists y \in T$  with  $\nabla r(y) = g$

## Idea

- $x_{t+1} = \pi_S^r(y_t) = \operatorname{argmin}_{x \in S} D_r(x || y_t)$
- Update  $\nabla r(y_t)$  using  $p_t$

# Two Algorithms

$$\pi_S^r(y) = \operatorname{argmin}_{x \in S} D_r(x||y)$$

## Algorithm #1

- $y_0 : \nabla r(y_0) = 0$
- For  $t \geq 1$ 
  - $x_t = \pi_S^r(y_t)$
  - $y_{t+1} : \nabla r(y_{t+1}) := \nabla r(y_t) - \eta p_t$

## FTRL / Dual Averaging

- $\Phi_T(x) \stackrel{\text{def}}{=} \eta \sum_{t \in [T]} p_t^\top x + r(x)$
- Let  $x_{T+1} = \operatorname{argmin}_{x \in S} \Phi_T(x)$
- $y_t : \nabla r(y_t) = -\eta \sum_{t \in [T]} p_t^\top x$

## Algorithm #2

- $y_0 = x_1$
- For  $t \geq 1$ 
  - $x_t = \pi_S^r(y_t)$
  - $y_{t+1} : \nabla r(y_{t+1}) := \nabla r(x_t) - \eta p_t$

## Mirror Descent

- $x_{t+1} = \operatorname{argmin}_{x \in S} \eta \cdot p_t^\top x + D_r(x||x_t)$
- $y_t : \nabla r(y_t) = -\eta p_t + \nabla r(x_{t-1})$

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# Motivating Problem

## Finite Sum Optimization

- $\min_{x \in \mathbb{R}^n} F(x) = \frac{1}{m} \sum_{i \in [m]} f_i(x)$
- Each  $f_i$  is convex and  $L$ -smooth
- $F$  is  $\mu$ -strongly convex
- Have gradient oracle for  $f_i$  and want  $\epsilon$ -optimal point

## Note

- Natural machine learning problem
  - $x$  = model,  $i$  = data point,  $f_i$  = loss of model on data point  $i$
- Many variants

## Baseline?

- Gradient Descent
  - $F$  is at most  $m \cdot \frac{1}{\mu} \cdot L$ -smooth
  - $O(m(L/\mu) \log \epsilon^{-1})$  queries
- AGD solves
  - $O(m\sqrt{L/\mu} \log \epsilon^{-1})$  queries
- Can we improve?

# Example

## Example Problem: Regression

- $\min_{x \in \mathbb{R}^n} F(x) = \frac{1}{2} \|Ax - b\|_2^2$
- $A$  has rows  $a_1, \dots, a_m \in \mathbb{R}^n$
- $f_i(x) = \frac{m}{2} (a_i^\top x - b)^2$
- $F(x) = \frac{1}{m} \sum_{i \in [m]} f_i(x)$
- Each  $f_i$  is  $\|a_i\|_2^2 m \leq L$  smooth

- $\min_{x \in \mathbb{R}^n} F(x) = \frac{1}{m} \sum_{i \in [m]} f_i(x)$
- Each  $f_i$  is convex and  $L$ -smooth
- $F$  is  $\mu$ -strongly convex
- Have gradient oracle for  $f_i$  and want  $\epsilon$ -optimal point

## Baseline

- $F$  is  $\mu = \lambda_{\min}(A^\top A)$ -strongly convex
- $F$  is  $\lambda_{\max}(A^\top A)$ -smooth
- $\lambda_{\max}(A^\top A) \leq \text{tr}(A^\top A) = \sum_{i \in [m]} \|a_i\|_2^2 \leq L$ 
  - Equal in the worst case
  - Can be better
    - $\lambda_{\max}(A^\top A) \geq \text{tr}(A^\top A)/d$
- GD:  $O(m(L/\mu) \log \epsilon^{-1})$  rows
- AGD:  $O(m\sqrt{L/\mu} \log \epsilon^{-1})$  rows
- Can we improve?

# Stochastic Gradient Descent (SGD)

## Problem

- Differentiable  $f: \mathbb{R}^n \rightarrow \mathbb{R}$
- Goal:  $\min_x f(x)$

## Algorithm

- For  $t \geq 1$ 
  - $g_t$  random s.t.  $\mathbb{E} g_t = \nabla f(x_t)$
  - $x_{t+1} = x_t + \eta_t g_t$

## Notes

- Many many variant
- Can have  $g_t \in \partial f(x_t)$
- Stochastic mirror descent
- Various  $\eta_t$  and  $x_t$  aggregation schemes
- Etc.
- **Today:** simple case of differentiable convex  $f$

# SGD

- $\min_{x \in \mathbb{R}^n} F(x) = \frac{1}{m} \sum_{i \in [m]} f_i(x)$
- Each  $f_i$  is convex and  $L$ -smooth
- $F$  is  $\mu$ -strongly convex
- Have gradient oracle for  $f_i$  and want  $\epsilon$ -optimal point
- $x_*$  minimizes and  $\epsilon_t \stackrel{\text{def}}{=} \mathbb{E}f(x_t) - f_*$

SGD Step: If  $x_{t+1} = x_t - \eta g_t$  and  $\mathbb{E}g_t = \nabla f(x_t)$

- $\epsilon_t \leq \frac{1}{2\eta} [\|x_t - x_*\|_2^2 - \mathbb{E}\|x_{t+1} - x_*\|_2^2] + \frac{\eta}{2} \mathbb{E}\|g_t\|_2^2$
- $\frac{1}{T} \sum_{t \in [T]} \epsilon_t \leq \frac{1}{T\eta\mu} \cdot \epsilon_1 + \frac{\eta}{2T} \sum_{t \in [T]} \mathbb{E}\|g_t\|_2^2$

Question  
How pick  $g_t$  ?

Proof:

- $\|x_{t+1} - x_*\|_2^2 = \|x_t - x_* - \eta g_t\|_2^2 = \|x_t - x_*\|_2^2 - 2\eta g_t^\top (x_t - x_*) + \eta^2 \|g_t\|_2^2$
- $\mathbb{E}g_t^\top (x_t - x_*) = \nabla f(x_t)^\top (x_t - x_*) = -\nabla f(x_t)^\top (x_* - x_t) \leq -[f(x_t) - f_*]$
- $\|x_{T+1} - x_*\|_2^2 \geq 0$
- $\|x_1 - x_*\|_2^2 \leq \frac{2}{\mu} \cdot \epsilon_1$

Note: will be a little informal (e.g. with expectations in slides, will be more formal in notes).

# Random Function

- $\min_{x \in \mathbb{R}^n} F(x) = \frac{1}{m} \sum_{i \in [m]} f_i(x)$
- Each  $f_i$  is convex and  $L$ -smooth
- $F$  is  $\mu$ -strongly convex
- Have gradient oracle for  $f_i$  and want  $\epsilon$ -optimal point
- $x_*$  minimizes and  $\epsilon_t \stackrel{\text{def}}{=} \mathbb{E}f(x_t) - f_*$

**SGD Step:** If  $x_{t+1} = x_t + \eta g_t$  and  $\mathbb{E}g_t = \nabla f(x_t)$

- $\frac{1}{T} \sum_{t \in [T]} \epsilon_t \leq \frac{1}{T\eta\mu} \cdot \epsilon_1 + \frac{\eta}{2T} \sum_{t \in [T]} \mathbb{E}\|g_t\|_2^2$
- **Idea:**  $g_t = \nabla f_{i_t}(x_t)$  for random  $i_t \in [m]$

## Note

Are prominent cases where  $\epsilon_i$  is 0, e.g.  $\epsilon_i = 0$   
(Kaczmarz)  $Ax = b$  where  $\exists x_*$  s.t.  $Ax_* = b$

## Analysis

- $\mathbb{E}\|g_t\|_2^2 = \mathbb{E}_{i_k} \|\nabla f_{i_k}(x_k) - \nabla f_{i_k}(x_*) + \nabla f_{i_k}(x_*)\|_2^2$
- $\mathbb{E}\|g_t\|_2^2 \leq 2 \cdot \mathbb{E}_{i_k} \|\nabla f_{i_k}(x_k) - \nabla f_{i_k}(x_*)\|_2^2 + 2 \cdot \mathbb{E}_{i_k} \|\nabla f_{i_k}(x_*)\|_2^2$
- $\|\nabla f_{i_k}(x_k) - \nabla f_{i_k}(x_*)\|_2^2 \leq 2L \cdot \left[ f_{i_k}(x_k) - [f_{i_k}(x_*) - \nabla f_{i_k}(x_*)^\top (x_k - x_*)] \right]$
- $\mathbb{E}_{i_k} \|\nabla f_{i_k}(x_k) - \nabla f_{i_k}(x_*)\|_2^2 \leq 2L \cdot [F(x_k) - F_*] \leq 2L \cdot \epsilon_k$

## Problem

Can be non-zero and large!



## Example

$f_i(x) = (a_i^\top x - b_i)^2$   
and  $b_i = a_i^\top x_* + \epsilon_i$  where  
 $\mathbb{E}\epsilon_i = 0$  and  $m \rightarrow \infty$

**Great!**  
Relative error

# Decreasing “Variance”

- $\min_{x \in \mathbb{R}^n} F(x) = \frac{1}{m} \sum_{i \in [m]} f_i(x)$
- Each  $f_i$  is convex and  $L$ -smooth
- $F$  is  $\mu$ -strongly convex
- Have gradient oracle for  $f_i$  and want  $\epsilon$ -optimal point
- $x_*$  minimizes and  $\epsilon_t \stackrel{\text{def}}{=} \mathbb{E}f(x_t) - f_*$

## Problem

- Want  $\mathbb{E}\|g_t\|_2^2$  small relative to the  $\epsilon_i$
- Difficulty: even  $\mathbb{E}_{i_k} \|\nabla f_{i_k}(x_*)\|_2^2 \neq 0$

Idea: add mean  $\vec{0}$  vector to improve variance

- $g_t = \nabla f_{i_t}(x_t) - \nabla f_{i_t}(x_*)$    **Pro:**  $\mathbb{E}\|g_t\|_2^2 = 0$  when  $x_1 = x_*$    **Con:** need know  $x_*$
- $g_t = \nabla f_{i_t}(x_t) - \nabla f_{i_t}(x_1)$    **Pro:**  $\mathbb{E}\|g_t\|_2^2 = 0$  when  $x_1 = x_*$    **Con:** biased!
- $g_t = \nabla f_{i_t}(x_t) - \nabla f_{i_t}(x_1) + \nabla F(x_1)$

$m$ -queries for initial  $g_1$  then 1 query for each additional  $g_t$

**Stochastic Variance Reduced Gradient (SVRG)**  
Johnson Zhang 13

# SVRG Variance

- $\min_{x \in \mathbb{R}^n} F(x) = \frac{1}{m} \sum_{i \in [m]} f_i(x)$
- Each  $f_i$  is convex and  $L$ -smooth
- $F$  is  $\mu$ -strongly convex
- Have gradient oracle for  $f_i$  and want  $\epsilon$ -optimal point
- $x_*$  minimizes and  $\epsilon_t \stackrel{\text{def}}{=} \mathbb{E}f(x_t) - f_*$

**SVRG Step:**  $g_t = \nabla f_{i_t}(x_t) - \nabla f_{i_t}(x_1) + \nabla F(x_1)$  for random  $i_t \in [m]$

- $\mathbb{E}\|g_t\|_2^2 = \mathbb{E}\|\nabla f_{i_t}(x_t) - \nabla f_{i_t}(x_*) - [\nabla f_{i_t}(x_1) - \nabla f_{i_t}(x_*)] + \nabla F(x_1)\|_2^2$
- $\mathbb{E}\|g_t\|_2^2 \leq 2 \cdot \mathbb{E}\|\nabla f_{i_t}(x_t) - \nabla f_{i_t}(x_*)\|_2^2$ 

$$+ 2 \cdot \mathbb{E}\|\nabla f_{i_t}(x_1) - \nabla f_{i_t}(x_*) - \nabla F(x_1)\|_2^2$$
- $\mathbb{E}\|v - \mathbb{E}v\|_2^2 \leq \mathbb{E}\|v\|_2^2 \Rightarrow$
- $\mathbb{E}\|g_t\|_2^2 \leq 2 \cdot [\mathbb{E}\|\nabla f_{i_t}(x_t) - \nabla f_{i_t}(x_*)\|_2^2 + \mathbb{E}\|\nabla f_{i_t}(x_1) - \nabla f_{i_t}(x_*)\|_2^2]$
- $\mathbb{E}\|g_t\|_2^2 \leq 4L[\epsilon_t + \epsilon_1]$

# SVRG

- $\min_{x \in \mathbb{R}^n} F(x) = \frac{1}{m} \sum_{i \in [m]} f_i(x)$
- Each  $f_i$  is convex and  $L$ -smooth
- $F$  is  $\mu$ -strongly convex
- Have gradient oracle for  $f_i$  and want  $\epsilon$ -optimal point
- $x_*$  minimizes and  $\epsilon_t \stackrel{\text{def}}{=} \mathbb{E}f(x_t) - f_*$

**SVRG:**  $x_{t+1} = x_t + \eta g_t$  and  $g_t = \nabla f_{i_t}(x_t) - \nabla f_{i_t}(x_1) + \nabla F(x_1)$  for  $i_t$  uniform

- $\Rightarrow \frac{1}{T} \sum_{t \in [T]} \epsilon_t \leq \frac{1}{T\eta\mu} \cdot \epsilon_1 + \frac{\eta}{2T} \sum_{t \in [T]} \mathbb{E}\|g_t\|_2^2$
- $\Rightarrow \frac{1}{T} \sum_{t \in [T]} \epsilon_t \leq \frac{1}{T} \left[ \frac{1}{\eta\mu} \cdot \epsilon_1 + \frac{\eta}{2} \cdot 4L \sum_{t \in [T]} [\epsilon_t + \epsilon_1] \right]$
- $\Rightarrow \frac{1-2\eta L}{T} \sum_{t \in [T]} \epsilon_t \leq \frac{1}{\eta\mu T} \cdot \epsilon_1 + 2L\eta\epsilon_1$
- $\left( \eta = \frac{1}{8L} \right) \Rightarrow \frac{1}{T} \sum_{t \in [T]} \epsilon_T \leq \frac{16L}{\mu T} \cdot \epsilon_1 + \frac{1}{2} \epsilon_1$
- $\left( T = \frac{64L}{\mu} \right) \Rightarrow \frac{1}{T} \sum_{t \in [T]} \epsilon_T \leq \frac{3}{4} \epsilon_1$

## Query Complexity?

- Can halve error in expectation with  $m + T = O(m + L/\mu)$  queries.
- Repeating yields expected  $\epsilon$ -optimal in  $O((m + L/\mu) \log \epsilon^{-1})$  queries!
- GD:  $O(m \cdot (L/\mu) \cdot \log \epsilon^{-1})$

# State-of-the-art?

- $\min_{x \in \mathbb{R}^n} F(x) = \frac{1}{m} \sum_{i \in [m]} f_i(x)$
- Each  $f_i$  is convex and  $L$ -smooth
- $F$  is  $\mu$ -strongly convex
- Have gradient oracle for  $f_i$  and want  $\epsilon$ -optimal point
- $x_*$  minimizes and  $\epsilon_t \stackrel{\text{def}}{=} \mathbb{E}f(x_t) - f_*$

- SVRG:  $O((m + L/\mu) \log \epsilon^{-1})$  queries
  - Are other ways to achieve
- GD:  $O(m \cdot (L/\mu) \cdot \log \epsilon^{-1})$  queries
- AGD:  $O(m \cdot (L/\mu) \cdot \log \epsilon^{-1})$  queries
- Accelerated proximal point / Catalyst
  - $\tilde{O}(m + \sqrt{mL/\mu}) \log \epsilon^{-1}$
  - Many generalization!
  - Many more applications!
    - MDPs [SWWY18, SWWYY18]
    - Submodular [CLSW17, HRRS19, ALS19]
    - Minimax [CJST19]
    - Regression, non-convex, ...
- Katyusha
  - $O(m + \sqrt{mL/\mu}) \log \epsilon^{-1}$

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